

Evaluate Winding Numbers through Cauchy Indices

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Abstract

In complex analysis, the winding number measures the number of times a path (counterclockwise) winds around a point, while the Cauchy index can approximate how the path winds. This entry provides a formalisation of the Cauchy index, which is then shown to be related to the winding number. In addition, this entry also offers a tactic that enables users to evaluate the winding number by calculating Cauchy indices. The connection between the winding number and the Cauchy index can be found in the literature [1] [2, Chapter 11].

1 Some useful lemmas in topology

theory *Missing-Topology* **imports** *HOL-Analysis.Multivariate-Analysis*
begin

1.1 Misc

lemma *open-times-image*:

fixes $S::'a::\text{real-normed-field set}$

assumes *open S c \neq 0*

shows *open (($*$) c) ' S*

<proof>

lemma *image-linear-greaterThan*:

fixes $x::'a::\text{linordered-field}$

assumes *c \neq 0*

shows $((\lambda x. c*x+b) ' \{x<..\}) = (\text{if } c>0 \text{ then } \{c*x+b <..\} \text{ else } \{..< c*x+b\})$

<proof>

lemma *image-linear-lessThan*:

fixes $x::'a::\text{linordered-field}$

assumes *c \neq 0*

shows $((\lambda x. c*x+b) ' \{..<x\}) = (\text{if } c>0 \text{ then } \{..<c*x+b\} \text{ else } \{c*x+b<..\})$

<proof>

lemma *continuous-on-neq-split*:

fixes $f :: 'a::\text{linear-continuum-topology} \Rightarrow 'b::\text{linorder-topology}$

assumes $\forall x \in s. f x \neq y$ *continuous-on s f connected s*
shows $(\forall x \in s. f x > y) \vee (\forall x \in s. f x < y)$
 <proof>

lemma
fixes $f::'a::\text{linorder-topology} \Rightarrow 'b::\text{topological-space}$
assumes *continuous-on {a..b} f a < b*
shows *continuous-on-at-left:continuous (at-left b) f*
and *continuous-on-at-right:continuous (at-right a) f*
 <proof>

1.2 More about *eventually*

lemma *eventually-comp-filtermap*:
eventually (P o f) F \longleftrightarrow eventually P (filtermap f F)
 <proof>

lemma *eventually-uminus-at-top-at-bot*:
fixes $P::'a::\{\text{ordered-ab-group-add,linorder}\} \Rightarrow \text{bool}$
shows *eventually (P o uminus) at-bot \longleftrightarrow eventually P at-top*
eventually (P o uminus) at-top \longleftrightarrow eventually P at-bot
 <proof>

lemma *eventually-at-infinityI*:
fixes $P::'a::\text{real-normed-vector} \Rightarrow \text{bool}$
assumes $\bigwedge x. c \leq \text{norm } x \implies P x$
shows *eventually P at-infinity*
 <proof>

lemma *eventually-at-bot-linorderI*:
fixes $c::'a::\text{linorder}$
assumes $\bigwedge x. x \leq c \implies P x$
shows *eventually P at-bot*
 <proof>

lemma *eventually-times-inverse-1*:
fixes $f::'a \Rightarrow 'b::\{\text{field,t2-space}\}$
assumes $(f \longrightarrow c) F c \neq 0$
shows $\forall_F x \text{ in } F. \text{inverse } (f x) * f x = 1$
 <proof>

1.3 More about *filtermap*

lemma *filtermap-linear-at-within*:
assumes *bij f and cont: isCont f a and open-map: $\bigwedge S. \text{open } S \implies \text{open } (f'S)$*
shows *filtermap f (at a within S) = at (f a) within f'S*
 <proof>

lemma *filtermap-at-bot-linear-eq*:
fixes $c::'a::\text{linordered-field}$

assumes $c \neq 0$
shows $\text{filtermap } (\lambda x. x * c + b) \text{ at-bot} = (\text{if } c > 0 \text{ then at-bot else at-top})$
 $\langle \text{proof} \rangle$

lemma *filtermap-linear-at-left*:
fixes $c :: 'a :: \{\text{linordered-field}, \text{linorder-topology}, \text{real-normed-field}\}$
assumes $c \neq 0$
shows $\text{filtermap } (\lambda x. c * x + b) (\text{at-left } x) = (\text{if } c > 0 \text{ then at-left } (c * x + b) \text{ else at-right } (c * x + b))$
 $\langle \text{proof} \rangle$

lemma *filtermap-linear-at-right*:
fixes $c :: 'a :: \{\text{linordered-field}, \text{linorder-topology}, \text{real-normed-field}\}$
assumes $c \neq 0$
shows $\text{filtermap } (\lambda x. c * x + b) (\text{at-right } x) = (\text{if } c > 0 \text{ then at-right } (c * x + b) \text{ else at-left } (c * x + b))$
 $\langle \text{proof} \rangle$

lemma *filtermap-at-top-linear-eq*:
fixes $c :: 'a :: \text{linordered-field}$
assumes $c \neq 0$
shows $\text{filtermap } (\lambda x. x * c + b) \text{ at-top} = (\text{if } c > 0 \text{ then at-top else at-bot})$
 $\langle \text{proof} \rangle$

lemma *filtermap-nhds-open-map*:
assumes $\text{cont}: \text{isCont } f \ a$
and $\text{open-map}: \bigwedge S. \text{open } S \implies \text{open } (f \cdot S)$
shows $\text{filtermap } f (\text{nhds } a) = \text{nhds } (f \ a)$
 $\langle \text{proof} \rangle$

1.4 More about *filterlim*

lemma *filterlim-at-infinity-times*:
fixes $f :: 'a \Rightarrow 'b :: \text{real-normed-field}$
assumes $\text{filterlim } f \text{ at-infinity } F \ \text{filterlim } g \text{ at-infinity } F$
shows $\text{filterlim } (\lambda x. f \ x * g \ x) \text{ at-infinity } F$
 $\langle \text{proof} \rangle$

lemma *filterlim-at-top-at-bot[elim]*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{unbounded-dense-linorder}, \text{order-topology}\}$ **and** $F :: 'a \text{ filter}$
assumes $\text{top}: \text{filterlim } f \text{ at-top } F$ **and** $\text{bot}: \text{filterlim } f \text{ at-bot } F$ **and** $F \neq \text{bot}$
shows False
 $\langle \text{proof} \rangle$

lemma *filterlim-at-top-nhds[elim]*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{unbounded-dense-linorder}, \text{order-topology}\}$ **and** $F :: 'a \text{ filter}$
assumes $\text{top}: \text{filterlim } f \text{ at-top } F$ **and** $\text{tendsto}: (f \longrightarrow c) \ F$ **and** $F \neq \text{bot}$
shows False

<proof>

lemma *filterlim-at-bot-nhds[elim]*:

fixes $f::'a \Rightarrow 'b::\{\text{unbounded-dense-linorder}, \text{order-topology}\}$ **and** $F::'a \text{ filter}$
assumes $\text{top:filterlim } f \text{ at-bot } F$ **and** $\text{tendsto: } (f \longrightarrow c) F$ **and** $F \neq \text{bot}$
shows *False*

<proof>

lemma *filterlim-at-top-linear-iff*:

fixes $f::'a::\text{linordered-field} \Rightarrow 'b$
assumes $c \neq 0$
shows $(\text{LIM } x \text{ at-top. } f (x * c + b) :> F2) \iff (\text{if } c > 0 \text{ then } (\text{LIM } x \text{ at-top. } f x :> F2) \text{ else } (\text{LIM } x \text{ at-bot. } f x :> F2))$

<proof>

lemma *filterlim-at-bot-linear-iff*:

fixes $f::'a::\text{linordered-field} \Rightarrow 'b$
assumes $c \neq 0$
shows $(\text{LIM } x \text{ at-bot. } f (x * c + b) :> F2) \iff (\text{if } c > 0 \text{ then } (\text{LIM } x \text{ at-bot. } f x :> F2) \text{ else } (\text{LIM } x \text{ at-top. } f x :> F2))$

<proof>

lemma *filterlim-tendsto-add-at-top-iff*:

assumes $f: (f \longrightarrow c) F$
shows $(\text{LIM } x F. (f x + g x :: \text{real}) :> \text{at-top}) \iff (\text{LIM } x F. g x :> \text{at-top})$

<proof>

lemma *filterlim-tendsto-add-at-bot-iff*:

fixes $c::\text{real}$
assumes $f: (f \longrightarrow c) F$
shows $(\text{LIM } x F. f x + g x :> \text{at-bot}) \iff (\text{LIM } x F. g x :> \text{at-bot})$

<proof>

lemma *tendsto-inverse-0-at-infinity*:

$\text{LIM } x F. f x :> \text{at-infinity} \implies ((\lambda x. \text{inverse } (f x) :: \text{real}) \longrightarrow 0) F$

<proof>

lemma *filterlim-at-infinity-divide-iff*:

fixes $f::'a \Rightarrow 'b::\text{real-normed-field}$
assumes $(f \longrightarrow c) F$ $c \neq 0$
shows $(\text{LIM } x F. f x / g x :> \text{at-infinity}) \iff (\text{LIM } x F. g x :> \text{at } 0)$

<proof>

lemma *filterlim-tendsto-pos-mult-at-top-iff*:

fixes $f::'a \Rightarrow \text{real}$

assumes $(f \longrightarrow c) F$ **and** $0 < c$
shows $(LIM\ x\ F.\ (f\ x\ * \ g\ x)\ :>\ at\ top) \longleftrightarrow (LIM\ x\ F.\ g\ x\ :>\ at\ top)$
 $\langle proof \rangle$

lemma *filterlim-tendsto-pos-mult-at-bot-iff*:
fixes $c :: real$
assumes $(f \longrightarrow c) F$ $0 < c$
shows $(LIM\ x\ F.\ f\ x\ * \ g\ x\ :>\ at\ bot) \longleftrightarrow filterlim\ g\ at\ bot\ F$
 $\langle proof \rangle$

lemma *filterlim-tendsto-neg-mult-at-top-iff*:
fixes $f :: 'a \Rightarrow real$
assumes $(f \longrightarrow c) F$ **and** $c < 0$
shows $(LIM\ x\ F.\ (f\ x\ * \ g\ x)\ :>\ at\ top) \longleftrightarrow (LIM\ x\ F.\ g\ x\ :>\ at\ bot)$
 $\langle proof \rangle$

lemma *filterlim-tendsto-neg-mult-at-bot-iff*:
fixes $c :: real$
assumes $(f \longrightarrow c) F$ $0 > c$
shows $(LIM\ x\ F.\ f\ x\ * \ g\ x\ :>\ at\ bot) \longleftrightarrow filterlim\ g\ at\ top\ F$
 $\langle proof \rangle$

lemma *Lim-add*:
fixes $f\ g :: 'a \Rightarrow \{t2\text{-space}, topological\text{-monoid-add}\}$
assumes $\exists y.\ (f \longrightarrow y) F$ **and** $\exists y.\ (g \longrightarrow y) F$ **and** $F \neq bot$
shows $Lim\ F\ f\ +\ Lim\ F\ g = Lim\ F\ (\lambda x.\ f\ x\ +\ g\ x)$
 $\langle proof \rangle$

1.5 Isolate and discrete

definition (in *topological-space*) *isolate*:: $'a\ set \Rightarrow bool$ (**infixr** *isolate* 60)
where $x\ isolate\ S \longleftrightarrow (x \in S \wedge (\exists T.\ open\ T \wedge T \cap S = \{x\}))$

definition (in *topological-space*) *discrete*:: $'a\ set \Rightarrow bool$
where $discrete\ S \longleftrightarrow (\forall x \in S.\ x\ isolate\ S)$

definition (in *metric-space*) *uniform-discrete* :: $'a\ set \Rightarrow bool$ **where**
 $uniform\text{-discrete}\ S \longleftrightarrow (\exists e > 0.\ \forall x \in S.\ \forall y \in S.\ dist\ x\ y < e \longrightarrow x = y)$

lemma *uniformI1*:
assumes $e > 0 \wedge x\ y.\ [x \in S; y \in S; dist\ x\ y < e] \Longrightarrow x = y$
shows $uniform\text{-discrete}\ S$
 $\langle proof \rangle$

lemma *uniformI2*:
assumes $e > 0 \wedge x\ y.\ [x \in S; y \in S; x \neq y] \Longrightarrow dist\ x\ y \geq e$
shows $uniform\text{-discrete}\ S$
 $\langle proof \rangle$

lemma *isolate-islimpt-iff*: $(x \text{ isolate } S) \longleftrightarrow (\neg (x \text{ islimpt } S) \wedge x \in S)$
<proof>

lemma *isolate-dist-Ex-iff*:
fixes $x :: 'a :: \text{metric-space}$
shows $x \text{ isolate } S \longleftrightarrow (x \in S \wedge (\exists e > 0. \forall y \in S. \text{dist } x \ y < e \longrightarrow y = x))$
<proof>

lemma *discrete-empty[simp]*: *discrete* {}
<proof>

lemma *uniform-discrete-empty[simp]*: *uniform-discrete* {}
<proof>

lemma *isolate-insert*:
fixes $x :: 'a :: \text{t1-space}$
shows $x \text{ isolate } (\text{insert } a \ S) \longleftrightarrow x \text{ isolate } S \vee (x = a \wedge \neg (x \text{ islimpt } S))$
<proof>

lemma *uniform-discrete-imp-closed*:
uniform-discrete $S \implies \text{closed } S$
<proof>

lemma *uniform-discrete-imp-discrete*:
uniform-discrete $S \implies \text{discrete } S$
<proof>

lemma *isolate-subset*: $x \text{ isolate } S \implies T \subseteq S \implies x \in T \implies x \text{ isolate } T$
<proof>

lemma *discrete-subset[elim]*: *discrete* $S \implies T \subseteq S \implies \text{discrete } T$
<proof>

lemma *uniform-discrete-subset[elim]*: *uniform-discrete* $S \implies T \subseteq S \implies \text{uniform-discrete } T$
<proof>

lemma *continuous-on-discrete*: *discrete* $S \implies \text{continuous-on } S \ f$
<proof>

lemma *uniform-discrete-insert*:
fixes $S :: 'a :: \text{euclidean-space set}$
shows $\text{uniform-discrete } (\text{insert } a \ S) \longleftrightarrow \text{uniform-discrete } S$
<proof>

lemma *discrete-compact-finite-iff*:

fixes $S :: 'a::t1\text{-space set}$
shows $\text{discrete } S \wedge \text{compact } S \longleftrightarrow \text{finite } S$
 $\langle \text{proof} \rangle$

lemma *uniform-discrete-finite-iff*:
fixes $S :: 'a::\text{heine-borel set}$
shows $\text{uniform-discrete } S \wedge \text{bounded } S \longleftrightarrow \text{finite } S$
 $\langle \text{proof} \rangle$

lemma *uniform-discrete-image-scale*:
fixes $f :: 'a::\text{euclidean-space} \Rightarrow 'b::\text{euclidean-space}$
assumes $\text{uniform-discrete } S$ **and** $\text{dist}:\forall x\in S. \forall y\in S. \text{dist } x \ y = c * \text{dist } (f \ x) \ (f \ y)$
shows $\text{uniform-discrete } (f \ ' S)$
 $\langle \text{proof} \rangle$

end

2 Some useful lemmas in algebra

theory *Missing-Algebraic imports*
HOL-Computational-Algebra.Polynomial-Factorial
HOL-Computational-Algebra.Fundamental-Theorem-Algebra
HOL-Complex-Analysis.Complex-Analysis
Missing-Topology
Budan-Fourier.BF-Misc
begin

2.1 Misc

lemma *poly-holomorphic-on[simp]*:
 $(\text{poly } p) \text{ holomorphic-on } s$
 $\langle \text{proof} \rangle$

lemma *order-zorder*:
fixes $p::\text{complex poly}$ **and** $z::\text{complex}$
assumes $p \neq 0$
shows $\text{order } z \ p = \text{nat } (\text{zorder } (\text{poly } p) \ z)$
 $\langle \text{proof} \rangle$

lemma *pcompose-pCons-0*: $\text{pcompose } p \ [a:] = [:\text{poly } p \ a:]$
 $\langle \text{proof} \rangle$

lemma *pcompose-coeff-0*:
 $\text{coeff } (\text{pcompose } p \ q) \ 0 = \text{poly } p \ (\text{coeff } q \ 0)$
 $\langle \text{proof} \rangle$

lemma *poly-field-differentiable-at[simp]*:
poly p field-differentiable (at x within s)
 ⟨proof⟩

lemma *deriv-pderiv*:
 $deriv (poly p) = poly (pderiv p)$
 ⟨proof⟩

lemma *lead-coeff-map-poly-nz*:
assumes $f (lead-coeff p) \neq 0$ $f 0 = 0$
shows $lead-coeff (map-poly f p) = f (lead-coeff p)$
 ⟨proof⟩

lemma *filterlim-poly-at-infinity*:
fixes $p::'a::real-normed-field$ *poly*
assumes $degree p > 0$
shows *filterlim (poly p) at-infinity at-infinity*
 ⟨proof⟩

lemma *poly-divide-tendsto-aux*:
fixes $p::'a::real-normed-field$ *poly*
shows $((\lambda x. poly p x / x^{(degree p)}) \longrightarrow lead-coeff p)$ *at-infinity*
 ⟨proof⟩

lemma *filterlim-power-at-infinity*:
assumes $n \neq 0$
shows *filterlim* $(\lambda x::'a::real-normed-field. x^n)$ *at-infinity at-infinity*
 ⟨proof⟩

lemma *poly-divide-tendsto-0-at-infinity*:
fixes $p::'a::real-normed-field$ *poly*
assumes $degree p > degree q$
shows $((\lambda x. poly q x / poly p x) \longrightarrow 0)$ *at-infinity*
 ⟨proof⟩

lemma *lead-coeff-list-def*:
 $lead-coeff p = (if coeffs p = [] then 0 else last (coeffs p))$
 ⟨proof⟩

lemma *poly-linepath-comp*:
fixes $a::'a::\{real-normed-vector, comm-semiring-0, real-algebra-1\}$
shows $poly p \circ (linepath a b) = poly (p \circ_p [:a, b-a:]) \circ of-real$
 ⟨proof⟩

lemma *poly-eventually-not-zero*:
fixes $p::real$ *poly*
assumes $p \neq 0$
shows *eventually* $(\lambda x. poly p x \neq 0)$ *at-infinity*
 ⟨proof⟩

2.2 More about *degree*

lemma *map-poly-degree-eq*:
 assumes f (*lead-coeff* p) $\neq 0$
 shows $\text{degree } (\text{map-poly } f \ p) = \text{degree } p$
 $\langle \text{proof} \rangle$

lemma *map-poly-degree-less*:
 assumes f (*lead-coeff* p) = 0 $\text{degree } p \neq 0$
 shows $\text{degree } (\text{map-poly } f \ p) < \text{degree } p$
 $\langle \text{proof} \rangle$

lemma *map-poly-degree-leq[simp]*:
 shows $\text{degree } (\text{map-poly } f \ p) \leq \text{degree } p$
 $\langle \text{proof} \rangle$

2.3 roots / zeros of a univariate function

definition *roots-within::('a \Rightarrow 'b::zero) \Rightarrow 'a set \Rightarrow 'a set* **where**
 $\text{roots-within } f \ s = \{x \in s. f \ x = 0\}$

abbreviation *roots::('a \Rightarrow 'b::zero) \Rightarrow 'a set* **where**
 $\text{roots } f \equiv \text{roots-within } f \ \text{UNIV}$

2.4 The argument principle specialised to polynomials.

lemma *argument-principle-poly*:
 assumes $p \neq 0$ **and** *valid:valid-path* g **and** *loop: pathfinish* $g = \text{pathstart } g$
 and *no-proots:path-image* $g \subseteq - \text{proots } p$
 shows $\text{contour-integral } g \ (\lambda x. \text{deriv } (\text{poly } p) \ x / \text{poly } p \ x) = 2 * \text{of-real } \pi * i *$
 $(\sum_{x \in \text{proots } p. \text{winding-number } g \ x * \text{of-nat } (\text{order } x \ p))$
 $\langle \text{proof} \rangle$

end

3 Some useful lemmas about transcendental functions

theory *Missing-Transcendental* **imports**
 Missing-Topology
 Missing-Algebraic
begin

3.1 Misc

lemma *Im-Ln-eq-pi-half*:
 $z \neq 0 \implies (\text{Im}(\text{Ln } z) = \pi/2 \iff 0 < \text{Im}(z) \wedge \text{Re}(z) = 0)$
 $z \neq 0 \implies (\text{Im}(\text{Ln } z) = -\pi/2 \iff \text{Im}(z) < 0 \wedge \text{Re}(z) = 0)$
 $\langle \text{proof} \rangle$

lemma *Im-Ln-eq*:

assumes $z \neq 0$

shows $\text{Im} (\text{Ln } z) = (\text{if } \text{Re } z \neq 0 \text{ then}$
 $\text{if } \text{Re } z > 0 \text{ then}$
 $\arctan (\text{Im } z / \text{Re } z)$
 $\text{else if } \text{Im } z \geq 0 \text{ then}$
 $\arctan (\text{Im } z / \text{Re } z) + \pi$
 else
 $\arctan (\text{Im } z / \text{Re } z) - \pi$
 $\text{if } \text{Im } z > 0 \text{ then } \pi/2 \text{ else } -\pi/2)$

<proof>

lemma *exp-Arg2pi2pi-multivalued*:

assumes $\exp (i * \text{of-real } x) = z$

shows $\exists k :: \text{int. } x = \text{Arg2pi } z + 2 * k * \pi$

<proof>

lemma *cos-eq-neg-periodic-intro*:

assumes $x - y = 2 * (\text{of-int } k) * \pi + \pi \vee x + y = 2 * (\text{of-int } k) * \pi + \pi$

shows $\cos x = - \cos y$ *<proof>*

lemma *cos-eq-periodic-intro*:

assumes $x - y = 2 * (\text{of-int } k) * \pi \vee x + y = 2 * (\text{of-int } k) * \pi$

shows $\cos x = \cos y$ *<proof>*

lemma *sin-tan-half*: $\sin (2 * x) = 2 * \tan x / (1 + (\tan x)^2)$

<proof>

lemma *cos-tan-half*: $\cos x \neq 0 \implies \cos (2 * x) = (1 - (\tan x)^2) / (1 + (\tan x)^2)$

<proof>

lemma *tan-eq-arctan-Ex*:

shows $\tan x = y \iff (\exists k :: \text{int. } x = \arctan y + k * \pi \vee (x = \pi/2 + k * \pi \wedge y = 0))$

<proof>

lemma *arccos-unique*:

assumes $0 \leq x$

and $x \leq \pi$

and $\cos x = y$

shows $\arccos y = x$

<proof>

lemma *cos-eq-arccos-Ex*:

$\cos x = y \iff -1 \leq y \wedge y \leq 1 \wedge (\exists k :: \text{int. } x = \arccos y + 2 * k * \pi \vee x = - \arccos y + 2 * k * \pi)$

<proof>

lemma *uniform-discrete-tan-eq*:
uniform-discrete { $x::\text{real}. \tan x = y$ }
 ⟨*proof*⟩

lemma *get-norm-value*:
fixes $a::'a::\{\text{floor-ceiling}\}$
assumes $pp>0$
obtains $k::\text{int}$ **and** $a1$ **where** $a=(\text{of-int } k)*pp+a1$ $a0\leq a1$ $a1<a0+pp$
 ⟨*proof*⟩

lemma *filtermap-tan-at-right*:
fixes $a::\text{real}$
assumes $\cos a\neq 0$
shows *filtermap tan* (*at-right* a) = *at-right* ($\tan a$)
 ⟨*proof*⟩

lemma *filtermap-tan-at-left*:
fixes $a::\text{real}$
assumes $\cos a\neq 0$
shows *filtermap tan* (*at-left* a) = *at-left* ($\tan a$)
 ⟨*proof*⟩

lemma *cos-zero-iff-int2*:
fixes $x::\text{real}$
shows $\cos x = 0 \iff (\exists n::\text{int}. x = n * \text{pi} + \text{pi}/2)$
 ⟨*proof*⟩

lemma *filtermap-tan-at-right-inf*:
fixes $a::\text{real}$
assumes $\cos a=0$
shows *filtermap tan* (*at-right* a) = *at-bot*
 ⟨*proof*⟩

lemma *filtermap-tan-at-left-inf*:
fixes $a::\text{real}$
assumes $\cos a=0$
shows *filtermap tan* (*at-left* a) = *at-top*
 ⟨*proof*⟩

3.2 Periodic set

definition *periodic-set*:: *real set* \Rightarrow *real* \Rightarrow *bool* **where**
periodic-set $S \delta \iff (\exists B. \text{finite } B \wedge (\forall x\in S. \exists b\in B. \exists k::\text{int}. x = b + k * \delta))$

lemma *periodic-set-multiple*:
assumes $k\neq 0$
shows *periodic-set* $S \delta \iff \text{periodic-set } S (\text{of-int } k*\delta)$

<proof>

lemma *periodic-set-empty[simp]*: *periodic-set* {} δ
<proof>

lemma *periodic-set-finite*:
 assumes *finite* S
 shows *periodic-set* S δ
<proof>

lemma *periodic-set-subset[elim]*:
 assumes *periodic-set* S δ $T \subseteq S$
 shows *periodic-set* T δ
<proof>

lemma *periodic-set-union*:
 assumes *periodic-set* S δ *periodic-set* T δ
 shows *periodic-set* $(S \cup T)$ δ
<proof>

lemma *periodic-imp-uniform-discrete*:
 assumes *periodic-set* S δ
 shows *uniform-discrete* S
<proof>

lemma *periodic-set-tan-linear*:
 assumes $a \neq 0$ $c \neq 0$
 shows *periodic-set* (*roots* $(\lambda x. a * \tan (x/c) + b)$) $(c * \pi)$
<proof>

lemma *periodic-set-cos-linear*:
 assumes $a \neq 0$ $c \neq 0$
 shows *periodic-set* (*roots* $(\lambda x. a * \cos (x/c) + b)$) $(2 * c * \pi)$
<proof>

lemma *periodic-set-tan-poly*:
 assumes $p \neq 0$ $c \neq 0$
 shows *periodic-set* (*roots* $(\lambda x. \text{poly } p (\tan (x/c)))$) $(c * \pi)$
<proof>

lemma *periodic-set-sin-cos-linear*:
 fixes a b $c :: \text{real}$
 assumes $a \neq 0 \vee b \neq 0 \vee c \neq 0$
 shows *periodic-set* (*roots* $(\lambda x. a * \cos x + b * \sin x + c)$) $(4 * \pi)$
<proof>

end

4 Some useful lemmas in analysis

```
theory Missing-Analysis
imports HOL-Complex-Analysis.Complex-Analysis
begin
```

4.1 More about paths

```
lemma pathfinish-offset[simp]:
  pathfinish ( $\lambda t. g\ t - z$ ) = pathfinish  $g - z$ 
   $\langle$ proof $\rangle$ 
```

```
lemma pathstart-offset[simp]:
  pathstart ( $\lambda t. g\ t - z$ ) = pathstart  $g - z$ 
   $\langle$ proof $\rangle$ 
```

```
lemma pathimage-offset[simp]:
  fixes  $g :: - \Rightarrow 'b::\text{topological-group-add}$ 
  shows  $p \in \text{path-image } (\lambda t. g\ t - z) \longleftrightarrow p+z \in \text{path-image } g$ 
   $\langle$ proof $\rangle$ 
```

```
lemma path-offset[simp]:
  fixes  $g :: - \Rightarrow 'b::\text{topological-group-add}$ 
  shows path ( $\lambda t. g\ t - z$ )  $\longleftrightarrow$  path  $g$ 
   $\langle$ proof $\rangle$ 
```

```
lemma not-on-circlepathI:
  assumes  $\text{cmod } (z-z0) \neq |r|$ 
  shows  $z \notin \text{path-image } (\text{part-circlepath } z0\ r\ st\ tt)$ 
   $\langle$ proof $\rangle$ 
```

```
lemma circlepath-inj-on:
  assumes  $r > 0$ 
  shows inj-on (circlepath  $z\ r$ )  $\{0..<1\}$ 
   $\langle$ proof $\rangle$ 
```

4.2 More lemmas related to winding-number

```
lemma winding-number-comp:
  assumes open  $s\ f$  holomorphic-on  $s$  path-image  $\gamma \subseteq s$ 
  valid-path  $\gamma\ z \notin \text{path-image } (f \circ \gamma)$ 
  shows winding-number  $(f \circ \gamma)\ z = 1/(2*\text{pi}*i)* \text{contour-integral } \gamma (\lambda w. \text{deriv } f\ w / (f\ w - z))$ 
   $\langle$ proof $\rangle$ 
```

```
lemma winding-number-uminus-comp:
  assumes valid-path  $\gamma - z \notin \text{path-image } \gamma$ 
  shows winding-number  $(\text{uminus} \circ \gamma)\ z = \text{winding-number } \gamma (-z)$ 
   $\langle$ proof $\rangle$ 
```

lemma *winding-number-comp-linear*:
assumes $c \neq 0$ *valid-path* γ **and** *not-image*: $(z-b)/c \notin \text{path-image } \gamma$
shows *winding-number* $((\lambda x. c*x+b) \circ \gamma) z = \text{winding-number } \gamma ((z-b)/c)$ **(is**
 $?L = ?R)$
 $\langle \text{proof} \rangle$

end

5 Cauchy's index theorem

theory *Cauchy-Index-Theorem* **imports**
HOL-Complex-Analysis.Complex-Analysis
Sturm-Tarski.Sturm-Tarski
HOL-Computational-Algebra.Fundamental-Theorem-Algebra
Missing-Transcendental
Missing-Algebraic
Missing-Analysis
begin

This theory formalises Cauchy indices on the complex plane and relate them to winding numbers

5.1 Misc

lemma *atMostAtLeast-subset-convex*:
fixes $C :: \text{real set}$
assumes *convex* C
and $x \in C$ $y \in C$
shows $\{x .. y\} \subseteq C$
 $\langle \text{proof} \rangle$

lemma *arg-elim*:
 $f x \implies x = y \implies f y$
 $\langle \text{proof} \rangle$

lemma *arg-elim2*:
 $f x1 x2 \implies x1 = y1 \implies x2 = y2 \implies f y1 y2$
 $\langle \text{proof} \rangle$

lemma *arg-elim3*:
 $\llbracket f x1 x2 x3; x1 = y1; x2 = y2; x3 = y3 \rrbracket \implies f y1 y2 y3$
 $\langle \text{proof} \rangle$

lemma *IVT-strict*:
fixes $f :: 'a::\text{linear-continuum-topology} \Rightarrow 'b::\text{linorder-topology}$
assumes $(f a > y \wedge y > f b) \vee (f a < y \wedge y < f b)$ $a < b$ *continuous-on* $\{a .. b\}$ f
shows $\exists x. a < x \wedge x < b \wedge f x = y$
 $\langle \text{proof} \rangle$

lemma (in *dense-linorder*) *atLeastAtMost-subseteq-greaterThanLessThan-iff*:
 $\{a \dots b\} \subseteq \{c < \dots < d\} \longleftrightarrow (a \leq b \longrightarrow c < a \wedge b < d)$
 ⟨proof⟩

lemma *Re-winding-number-half-right*:
assumes $\forall p \in \text{path-image } \gamma. \text{Re } p \geq \text{Re } z$ **and** *valid-path* γ **and** $z \notin \text{path-image } \gamma$
shows $\text{Re}(\text{winding-number } \gamma \ z) = (\text{Im } (\text{Ln } (\text{pathfinish } \gamma - z)) - \text{Im } (\text{Ln } (\text{pathstart } \gamma - z))) / (2 * \pi)$
 ⟨proof⟩

lemma *Re-winding-number-half-upper*:
assumes $\forall p \in \text{path-image } \gamma. \text{Im } p \geq \text{Im } z$ **and** *valid-path* γ **and** $z \notin \text{path-image } \gamma$
shows $\text{Re}(\text{winding-number } \gamma \ z) = (\text{Im } (\text{Ln } (i * z - i * \text{pathfinish } \gamma)) - \text{Im } (\text{Ln } (i * z - i * \text{pathstart } \gamma))) / (2 * \pi)$
 ⟨proof⟩

lemma *Re-winding-number-half-lower*:
assumes $\forall p \in \text{path-image } \gamma. \text{Im } p \leq \text{Im } z$ **and** *valid-path* γ **and** $z \notin \text{path-image } \gamma$
shows $\text{Re}(\text{winding-number } \gamma \ z) = (\text{Im } (\text{Ln } (i * \text{pathfinish } \gamma - i * z)) - \text{Im } (\text{Ln } (i * \text{pathstart } \gamma - i * z))) / (2 * \pi)$
 ⟨proof⟩

lemma *Re-winding-number-half-left*:
assumes *neg-imag*: $\forall p \in \text{path-image } \gamma. \text{Re } p \leq \text{Re } z$ **and** *valid-path* γ **and** $z \notin \text{path-image } \gamma$
shows $\text{Re}(\text{winding-number } \gamma \ z) = (\text{Im } (\text{Ln } (z - \text{pathfinish } \gamma)) - \text{Im } (\text{Ln } (z - \text{pathstart } \gamma))) / (2 * \pi)$
 ⟨proof⟩

lemma *continuous-on-open-Collect-neg*:
fixes $f \ g :: 'a :: \text{topological-space} \Rightarrow 'b :: \text{t2-space}$
assumes f : *continuous-on* $S \ f$ **and** g : *continuous-on* $S \ g$ **and** *open* S
shows *open* $\{x \in S. f \ x \neq g \ x\}$
 ⟨proof⟩

5.2 Sign at a filter

definition *has-sgnx*: $(\text{real} \Rightarrow \text{real}) \Rightarrow \text{real} \Rightarrow \text{real filter} \Rightarrow \text{bool}$
 (infixr *has'-sgnx* 55) **where**
 ($f \ \text{has-sgnx} \ c$) $F = (\text{eventually } (\lambda x. \text{sgn}(f \ x) = c) \ F)$

definition *sgnx-able* (infixr *sgnx'-able* 55) **where**
 ($f \ \text{sgnx-able} \ F$) = $(\exists c. (f \ \text{has-sgnx} \ c) \ F)$

definition *sgnx* **where**
 $\text{sgnx} \ f \ F = (\text{SOME } c. (f \ \text{has-sgnx} \ c) \ F)$

lemma *has-sgnx-eq-rhs*: $(f \text{ has-sgnx } x) F \implies x = y \implies (f \text{ has-sgnx } y) F$
 ⟨proof⟩

named-theorems *sgnx-intros* introduction rules for *has-sgnx*
 ⟨ML⟩

lemma *sgnx-able-sgnx*: $f \text{ sgnx-able } F \implies (f \text{ has-sgnx } (\text{sgnx } f F)) F$
 ⟨proof⟩

lemma *has-sgnx-imp-sgnx-able*[*elim*]:
 $(f \text{ has-sgnx } c) F \implies f \text{ sgnx-able } F$
 ⟨proof⟩

lemma *has-sgnx-unique*:
 assumes $F \neq \text{bot}$ $(f \text{ has-sgnx } c1) F$ $(f \text{ has-sgnx } c2) F$
 shows $c1 = c2$
 ⟨proof⟩

lemma *has-sgnx-imp-sgnx*[*elim*]:
 $(f \text{ has-sgnx } c) F \implies F \neq \text{bot} \implies \text{sgnx } f F = c$
 ⟨proof⟩

lemma *has-sgnx-const*[*simp,sgnx-intros*]:
 $((\lambda \cdot. c) \text{ has-sgnx } \text{sgn } c) F$
 ⟨proof⟩

lemma *finite-sgnx-at-left-at-right*:
 assumes *finite* $\{t. f t = 0 \wedge a < t \wedge t < b\}$ *continuous-on* $(\{a < .. < b\} - s)$ *f finite s*
 and $x : x \in \{a < .. < b\}$
 shows $f \text{ sgnx-able } (\text{at-left } x) \text{ sgnx } f (\text{at-left } x) \neq 0$
 $f \text{ sgnx-able } (\text{at-right } x) \text{ sgnx } f (\text{at-right } x) \neq 0$
 ⟨proof⟩

lemma *sgnx-able-poly*[*simp*]:
 $(\text{poly } p) \text{ sgnx-able } (\text{at-right } a)$
 $(\text{poly } p) \text{ sgnx-able } (\text{at-left } a)$
 $(\text{poly } p) \text{ sgnx-able at-top}$
 $(\text{poly } p) \text{ sgnx-able at-bot}$
 ⟨proof⟩

lemma *has-sgnx-identity*[*intro,sgnx-intros*]:
 shows $x \geq 0 \implies ((\lambda x. x) \text{ has-sgnx } 1) (\text{at-right } x)$
 $x \leq 0 \implies ((\lambda x. x) \text{ has-sgnx } -1) (\text{at-left } x)$
 ⟨proof⟩

lemma *has-sgnx-divide*[*sgnx-intros*]:
 assumes $(f \text{ has-sgnx } c1) F$ $(g \text{ has-sgnx } c2) F$
 shows $((\lambda x. f x / g x) \text{ has-sgnx } c1 / c2) F$

<proof>

lemma *sgnx-able-divide*[*sgnx-intros*]:
 assumes f *sgnx-able* F g *sgnx-able* F
 shows $(\lambda x. f\ x / g\ x)$ *sgnx-able* F
<proof>

lemma *sgnx-divide*:
 assumes $F \neq \text{bot}$ f *sgnx-able* F g *sgnx-able* F
 shows *sgnx* $(\lambda x. f\ x / g\ x)$ $F = \text{sgnx } f\ F / \text{sgnx } g\ F$
<proof>

lemma *has-sgnx-times*[*sgnx-intros*]:
 assumes $(f$ *has-sgnx* $c1)$ F $(g$ *has-sgnx* $c2)$ F
 shows $((\lambda x. f\ x * g\ x)$ *has-sgnx* $c1 * c2)$ F
<proof>

lemma *sgnx-able-times*[*sgnx-intros*]:
 assumes f *sgnx-able* F g *sgnx-able* F
 shows $(\lambda x. f\ x * g\ x)$ *sgnx-able* F
<proof>

lemma *sgnx-times*:
 assumes $F \neq \text{bot}$ f *sgnx-able* F g *sgnx-able* F
 shows *sgnx* $(\lambda x. f\ x * g\ x)$ $F = \text{sgnx } f\ F * \text{sgnx } g\ F$
<proof>

lemma *tendsto-nonzero-has-sgnx*:
 assumes $(f \longrightarrow c)$ F $c \neq 0$
 shows $(f$ *has-sgnx* *sgn* $c)$ F
<proof>

lemma *tendsto-nonzero-sgnx*:
 assumes $(f \longrightarrow c)$ F $F \neq \text{bot}$ $c \neq 0$
 shows *sgnx* $f\ F = \text{sgn } c$
<proof>

lemma *filterlim-divide-at-bot-at-top-iff*:
 assumes $(f \longrightarrow c)$ F $c \neq 0$
 shows
 $(\text{LIM } x\ F. f\ x / g\ x \text{ :> } \text{at-bot}) \iff (g \longrightarrow 0)$ F
 $\wedge ((\lambda x. g\ x)$ *has-sgnx* $-\text{sgn } c)$ F
 $(\text{LIM } x\ F. f\ x / g\ x \text{ :> } \text{at-top}) \iff (g \longrightarrow 0)$ F
 $\wedge ((\lambda x. g\ x)$ *has-sgnx* *sgn* $c)$ F
<proof>

lemma *poly-sgnx-left-right*:

fixes $c a::\text{real}$ **and** $p::\text{real poly}$
assumes $p \neq 0$
shows $\text{sgnx } (\text{poly } p) (\text{at-left } a) = (\text{if even } (\text{order } a) p)$
 $\quad \text{then } \text{sgnx } (\text{poly } p) (\text{at-right } a)$
 $\quad \text{else } -\text{sgnx } (\text{poly } p) (\text{at-right } a)$
 $\langle \text{proof} \rangle$

lemma *poly-has-sgnx-left-right*:
fixes $c a::\text{real}$ **and** $p::\text{real poly}$
assumes $p \neq 0$
shows $(\text{poly } p \text{ has-sgnx } c) (\text{at-left } a) \longleftrightarrow (\text{if even } (\text{order } a) p)$
 $\quad \text{then } (\text{poly } p \text{ has-sgnx } c) (\text{at-right } a)$
 $\quad \text{else } (\text{poly } p \text{ has-sgnx } -c) (\text{at-right } a)$
 $\langle \text{proof} \rangle$

lemma *sign-r-pos-sgnx-iff*:
 $\text{sign-r-pos } p a \longleftrightarrow \text{sgnx } (\text{poly } p) (\text{at-right } a) > 0$
 $\langle \text{proof} \rangle$

lemma *sgnx-values*:
assumes $f \text{ sgnx-able } F F \neq \text{bot}$
shows $\text{sgnx } f F = -1 \vee \text{sgnx } f F = 0 \vee \text{sgnx } f F = 1$
 $\langle \text{proof} \rangle$

lemma *has-sgnx-poly-at-top*:
 $(\text{poly } p \text{ has-sgnx } \text{sgn-pos-inf } p) \text{ at-top}$
 $\langle \text{proof} \rangle$

lemma *has-sgnx-poly-at-bot*:
 $(\text{poly } p \text{ has-sgnx } \text{sgn-neg-inf } p) \text{ at-bot}$
 $\langle \text{proof} \rangle$

lemma *sgnx-poly-at-top*:
 $\text{sgnx } (\text{poly } p) \text{ at-top} = \text{sgn-pos-inf } p$
 $\langle \text{proof} \rangle$

lemma *sgnx-poly-at-bot*:
 $\text{sgnx } (\text{poly } p) \text{ at-bot} = \text{sgn-neg-inf } p$
 $\langle \text{proof} \rangle$

lemma *poly-has-sgnx-values*:
assumes $p \neq 0$
shows
 $(\text{poly } p \text{ has-sgnx } 1) (\text{at-left } a) \vee (\text{poly } p \text{ has-sgnx } -1) (\text{at-left } a)$
 $(\text{poly } p \text{ has-sgnx } 1) (\text{at-right } a) \vee (\text{poly } p \text{ has-sgnx } -1) (\text{at-right } a)$
 $(\text{poly } p \text{ has-sgnx } 1) \text{ at-top} \vee (\text{poly } p \text{ has-sgnx } -1) \text{ at-top}$
 $(\text{poly } p \text{ has-sgnx } 1) \text{ at-bot} \vee (\text{poly } p \text{ has-sgnx } -1) \text{ at-bot}$
 $\langle \text{proof} \rangle$

lemma *poly-sgnx-values*:

assumes $p \neq 0$

shows $\text{sgnx}(\text{poly } p)(\text{at-left } a) = 1 \vee \text{sgnx}(\text{poly } p)(\text{at-left } a) = -1$

$\text{sgnx}(\text{poly } p)(\text{at-right } a) = 1 \vee \text{sgnx}(\text{poly } p)(\text{at-right } a) = -1$

<proof>

lemma *has-sgnx-inverse*: $(f \text{ has-sgnx } c) F \longleftrightarrow ((\text{inverse } o f) \text{ has-sgnx } (\text{inverse } c)) F$

<proof>

lemma *has-sgnx-derivative-at-left*:

assumes $g\text{-deriv}(g \text{ has-field-derivative } c)(\text{at } x)$ **and** $g \ x=0$ **and** $c \neq 0$

shows $(g \text{ has-sgnx } - \text{sgn } c)(\text{at-left } x)$

<proof>

lemma *has-sgnx-derivative-at-right*:

assumes $g\text{-deriv}(g \text{ has-field-derivative } c)(\text{at } x)$ **and** $g \ x=0$ **and** $c \neq 0$

shows $(g \text{ has-sgnx } \text{sgn } c)(\text{at-right } x)$

<proof>

lemma *has-sgnx-split*:

$(f \text{ has-sgnx } c)(\text{at } x) \longleftrightarrow (f \text{ has-sgnx } c)(\text{at-left } x) \wedge (f \text{ has-sgnx } c)(\text{at-right } x)$

<proof>

lemma *sgnx-at-top-IVT*:

assumes $\text{sgnx}(\text{poly } p)(\text{at-right } a) \neq \text{sgnx}(\text{poly } p) \text{ at-top}$

shows $\exists x > a. \text{poly } p \ x=0$

<proof>

lemma *sgnx-at-left-at-right-IVT*:

assumes $\text{sgnx}(\text{poly } p)(\text{at-right } a) \neq \text{sgnx}(\text{poly } p)(\text{at-left } b) \ a < b$

shows $\exists x. a < x \wedge x < b \wedge \text{poly } p \ x=0$

<proof>

lemma *sgnx-at-bot-IVT*:

assumes $\text{sgnx}(\text{poly } p)(\text{at-left } a) \neq \text{sgnx}(\text{poly } p) \text{ at-bot}$

shows $\exists x < a. \text{poly } p \ x=0$

<proof>

lemma *sgnx-poly-nz*:

assumes $\text{poly } p \ x \neq 0$

shows $\text{sgnx}(\text{poly } p)(\text{at-left } x) = \text{sgn}(\text{poly } p \ x)$

$\text{sgnx}(\text{poly } p)(\text{at-right } x) = \text{sgn}(\text{poly } p \ x)$

<proof>

5.3 Finite predicate segments over an interval

inductive *finite-Psegments*::(*real* \Rightarrow *bool*) \Rightarrow *real* \Rightarrow *real* \Rightarrow *bool* for *P* where

emptyI: $a \geq b \implies \text{finite-Psegments } P \ a \ b$

insertI-1: $\llbracket s \in \{a..<b\}; s = a \vee P \ s; \forall t \in \{s<..
 $\implies \text{finite-Psegments } P \ a \ b$$

insertI-2: $\llbracket s \in \{a..<b\}; s = a \vee P \ s; (\forall t \in \{s<..
 $\implies \text{finite-Psegments } P \ a \ b$$

lemma *finite-Psegments-pos-linear*:

assumes *finite-Psegments* *P* ($b * lb + c$) ($b * ub + c$) **and** $b > 0$

shows *finite-Psegments* (*P* o ($\lambda t. b * t + c$)) *lb ub*

<proof>

lemma *finite-Psegments-congE*:

assumes *finite-Psegments* *Q lb ub*

$\bigwedge t. \llbracket lb < t; t < ub \rrbracket \implies Q \ t \longleftrightarrow P \ t$

shows *finite-Psegments* *P lb ub* *<proof>*

lemma *finite-Psegments-constI*:

assumes $\bigwedge t. \llbracket a < t; t < b \rrbracket \implies P \ t = c$

shows *finite-Psegments* *P a b*

<proof>

context

begin

private lemma *finite-Psegments-less-eq1*:

assumes *finite-Psegments* *P a c b* $b \leq c$

shows *finite-Psegments* *P a b* *<proof>* **lemma** *finite-Psegments-less-eq2*:

assumes *finite-Psegments* *P a c a* $a \leq b$

shows *finite-Psegments* *P b c* *<proof>*

lemma *finite-Psegments-included*:

assumes *finite-Psegments* *P a d a* $a \leq b$ $c \leq d$

shows *finite-Psegments* *P b c*

<proof>

end

lemma *finite-Psegments-combine*:

assumes *finite-Psegments* *P a b* *finite-Psegments* *P b c* $b \in \{a..c\}$ *closed* ($\{x. P \ x\} \cap \{a..c\}$)

shows *finite-Psegments* *P a c* *<proof>*

5.4 Finite segment intersection of a path with the imaginary axis

definition *finite-ReZ-segments*::(*real* \Rightarrow *complex*) \Rightarrow *complex* \Rightarrow *bool* where

finite-ReZ-segments *g z = finite-Psegments* ($\lambda t. \text{Re } (g \ t - z) = 0$) *0 1*

lemma *finite-ReZ-segments-joinpaths*:

assumes $g1$:*finite-ReZ-segments* $g1$ z **and** $g2$: *finite-ReZ-segments* $g2$ z **and**
 $path$ $g1$ $path$ $g2$ $pathfinish$ $g1=pathstart$ $g2$
shows *finite-ReZ-segments* $(g1+++g2)$ z

<proof>

lemma *finite-ReZ-segments-congE*:

assumes *finite-ReZ-segments* $p1$ $z1$
 $\wedge t. \llbracket 0 < t; t < 1 \rrbracket \implies Re(p1\ t - z1) = Re(p2\ t - z2)$
shows *finite-ReZ-segments* $p2$ $z2$

<proof>

lemma *finite-ReZ-segments-constI*:

assumes $\forall t. 0 < t \wedge t < 1 \longrightarrow g\ t = c$
shows *finite-ReZ-segments* g z

<proof>

lemma *finite-ReZ-segment-cases* [*consumes 1, case-names subEq subNEq, cases pred:finite-ReZ-segments*]:

assumes *finite-ReZ-segments* g z
and *subEq*: $(\wedge s. \llbracket s \in \{0..<1\}; s=0 \vee Re(g\ s) = Re\ z; \forall t \in \{s<..$

P)

and *subNEq*: $(\wedge s. \llbracket s \in \{0..<1\}; s=0 \vee Re(g\ s) = Re\ z; \forall t \in \{s<..$

P)

shows P

<proof>

lemma *finite-ReZ-segments-induct* [*case-names sub0 subEq subNEq, induct pred:finite-ReZ-segments*]:

assumes *finite-ReZ-segments* g z
assumes *sub0*: $\wedge g\ z. (P(subpath\ 0\ 0\ g)\ z)$
and *subEq*: $(\wedge s\ g\ z. \llbracket s \in \{0..<1\}; s=0 \vee Re(g\ s) = Re\ z; \forall t \in \{s<..
and *subNEq*: $(\wedge s\ g\ z. \llbracket s \in \{0..<1\}; s=0 \vee Re(g\ s) = Re\ z; \forall t \in \{s<..$$

shows $P\ g\ z$

<proof>

lemma *finite-ReZ-segments-shiftpah*:

assumes *finite-ReZ-segments* g z $s \in \{0..1\}$ $path$ g **and** *loop*: $pathfinish\ g = pathstart\ g$

shows *finite-ReZ-segments* $(shiftpath\ s\ g)$ z

<proof>

lemma *finite-imp-finite-ReZ-segments*:

assumes *finite* $\{t. Re(g\ t - z) = 0 \wedge 0 \leq t \wedge t \leq 1\}$

shows *finite-ReZ-segments* $g z$
 ⟨*proof*⟩

lemma *finite-ReZ-segments-poly-linepath*:
shows *finite-ReZ-segments* (*poly* $p o$ *linepath* $a b$) z
 ⟨*proof*⟩

lemma *part-circlepath-half-finite-inter*:
assumes $st \neq tt \ r \neq 0 \ c \neq 0$
shows *finite* $\{t. \text{part-circlepath } z0 \ r \ st \ tt \ t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$ (**is** *finite* ? T)
 ⟨*proof*⟩

lemma *linepath-half-finite-inter*:
assumes $a \cdot c \neq d \vee b \cdot c \neq d$
shows *finite* $\{t. \text{linepath } a \ b \ t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$ (**is** *finite* ? S)
 ⟨*proof*⟩

lemma *finite-half-joinpaths-inter*:
assumes *finite* $\{t. l1 \ t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$ *finite* $\{t. l2 \ t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$
shows *finite* $\{t. (l1 +++ l2) \ t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$
 ⟨*proof*⟩

lemma *finite-ReZ-segments-linepath*:
finite-ReZ-segments (*linepath* $a b$) z
 ⟨*proof*⟩

lemma *finite-ReZ-segments-part-circlepath*:
finite-ReZ-segments (*part-circlepath* $z0 \ r \ st \ tt$) z
 ⟨*proof*⟩

lemma *finite-ReZ-segments-poly-of-real*:
shows *finite-ReZ-segments* (*poly* $p o$ *of-real*) z
 ⟨*proof*⟩

lemma *finite-ReZ-segments-subpath*:
assumes *finite-ReZ-segments* $g z$
 $0 \leq u \ u \leq v \ v \leq 1$
shows *finite-ReZ-segments* (*subpath* $u \ v \ g$) z
 ⟨*proof*⟩

5.5 jump and jumpF

definition *jump*::(*real* \Rightarrow *real*) \Rightarrow *real* \Rightarrow *int* **where**

jump $f a =$ (if
 (*LIM* x (*at-left* a). $f x$ $:$ \Rightarrow *at-bot*) \wedge (*LIM* x (*at-right* a). $f x$ $:$ \Rightarrow *at-top*)
 then 1 else if
 (*LIM* x (*at-left* a). $f x$ $:$ \Rightarrow *at-top*) \wedge (*LIM* x (*at-right* a). $f x$ $:$ \Rightarrow *at-bot*)
 then -1 else 0)

definition $jumpF::(real \Rightarrow real) \Rightarrow real\ filter \Rightarrow real$ **where**
 $jumpF\ f\ F \equiv (if\ filterlim\ f\ at-top\ F\ then\ 1/2\ else$
 $if\ filterlim\ f\ at-bot\ F\ then\ -1/2\ else\ (0::real))$

lemma $jumpF-const[simp]$:
assumes $F \neq bot$
shows $jumpF\ (\lambda-. c)\ F = 0$
 $\langle proof \rangle$

lemma $jumpF-not-infinity$:
assumes $continuous\ F\ g\ F \neq bot$
shows $jumpF\ g\ F = 0$
 $\langle proof \rangle$

lemma $jumpF-linear-comp$:
assumes $c \neq 0$
shows
 $jumpF\ (f\ o\ (\lambda x. c*x+b))\ (at-left\ x) =$
 $(if\ c>0\ then\ jumpF\ f\ (at-left\ (c*x+b))\ else\ jumpF\ f\ (at-right\ (c*x+b)))$
(is ?case1)
 $jumpF\ (f\ o\ (\lambda x. c*x+b))\ (at-right\ x) =$
 $(if\ c>0\ then\ jumpF\ f\ (at-right\ (c*x+b))\ else\ jumpF\ f\ (at-left\ (c*x+b)))$
(is ?case2)
 $\langle proof \rangle$

lemma $jump-const[simp]:jump\ (\lambda-. c)\ a = 0$
 $\langle proof \rangle$

lemma $jump-not-infinity$:
 $isCont\ f\ a \implies jump\ f\ a = 0$
 $\langle proof \rangle$

lemma $jump-jump-poly-aux$:
assumes $p \neq 0\ coprime\ p\ q$
shows $jump\ (\lambda x. poly\ q\ x / poly\ p\ x)\ a = jump-poly\ q\ p\ a$
 $\langle proof \rangle$

lemma $jump-jumpF$:
assumes $cont:isCont\ (inverse\ o\ f)\ a$ **and**
 $sgnxl:(f\ has-sgnx\ l)\ (at-left\ a)$ **and** $sgnxr:(f\ has-sgnx\ r)\ (at-right\ a)$ **and**
 $l \neq 0\ r \neq 0$
shows $jump\ f\ a = jumpF\ f\ (at-right\ a) - jumpF\ f\ (at-left\ a)$
 $\langle proof \rangle$

lemma $jump-linear-comp$:
assumes $c \neq 0$
shows $jump\ (f\ o\ (\lambda x. c*x+b))\ x = (if\ c>0\ then\ jump\ f\ (c*x+b)\ else\ -jump\ f\ (c*x+b))$

<proof>

lemma *jump-divide-derivative*:

assumes *isCont* f g $x = 0$ $f x \neq 0$

and *g-deriv*:(*g has-field-derivative* c) (*at* x) **and** $c \neq 0$

shows *jump* $(\lambda t. f t / g t) x = (\text{if } \text{sgn } c = \text{sgn } (f x) \text{ then } 1 \text{ else } -1)$

<proof>

lemma *jump-jump-poly*: *jump* $(\lambda x. \text{poly } q x / \text{poly } p x) a = \text{jump-poly } q p a$

<proof>

lemma *jump-Im-divide-Re-0*:

assumes *path* g *Re* $(g x) \neq 0$ $0 < x < 1$

shows *jump* $(\lambda t. \text{Im } (g t) / \text{Re } (g t)) x = 0$

<proof>

lemma *jumpF-im-divide-Re-0*:

assumes *path* g *Re* $(g x) \neq 0$

shows $\llbracket 0 \leq x; x < 1 \rrbracket \implies \text{jumpF } (\lambda t. \text{Im } (g t) / \text{Re } (g t)) (\text{at-right } x) = 0$

$\llbracket 0 < x; x \leq 1 \rrbracket \implies \text{jumpF } (\lambda t. \text{Im } (g t) / \text{Re } (g t)) (\text{at-left } x) = 0$

<proof>

lemma *jump-cong*:

assumes $x = y$ **and** *eventually* $(\lambda x. f x = g x)$ (*at* x)

shows *jump* $f x = \text{jump } g y$

<proof>

lemma *jumpF-cong*:

assumes $F = G$ **and** *eventually* $(\lambda x. f x = g x)$ F

shows *jumpF* $f F = \text{jumpF } g G$

<proof>

lemma *jump-at-left-at-right-eq*:

assumes *isCont* f x **and** $f x \neq 0$ **and** *sgnx-eq*:*sgnx* g (*at-left* x) = *sgnx* g (*at-right* x)

shows *jump* $(\lambda t. f t / g t) x = 0$

<proof>

lemma *jumpF-pos-has-sgnx*:

assumes *jumpF* $f F > 0$

shows (*f has-sgnx* 1) F

<proof>

lemma *jumpF-neg-has-sgnx*:

assumes *jumpF* $f F < 0$

shows (*f has-sgnx* -1) F

<proof>

lemma *jumpF-IVT*:

fixes $f::\text{real} \Rightarrow \text{real}$ **and** $a\ b::\text{real}$

defines $\text{right} \equiv (\lambda(R::\text{real} \Rightarrow \text{real} \Rightarrow \text{bool}). R (\text{jumpF } f (\text{at-right } a)) \ 0$
 $\quad \vee (\text{continuous } (\text{at-right } a) f \wedge R (f a) \ 0))$

and

$\text{left} \equiv (\lambda(R::\text{real} \Rightarrow \text{real} \Rightarrow \text{bool}). R (\text{jumpF } f (\text{at-left } b)) \ 0$
 $\quad \vee (\text{continuous } (\text{at-left } b) f \wedge R (f b) \ 0))$

assumes $a < b$ **and** $\text{cont}::\text{continuous-on } \{a <..<b\}$ f **and**

$\text{right-left}::\text{right greater} \wedge \text{left less} \vee \text{right less} \wedge \text{left greater}$

shows $\exists x. a < x \wedge x < b \wedge f x = 0$

$\langle \text{proof} \rangle$

lemma *jumpF-eventually-const*:

assumes $\text{eventually } (\lambda x. f x = c)$ F $F \neq \text{bot}$

shows $\text{jumpF } f F = 0$

$\langle \text{proof} \rangle$

lemma *jumpF-tan-comp*:

$\text{jumpF } (f \circ \text{tan}) (\text{at-right } x) = (\text{if } \cos x = 0$
 $\quad \text{then } \text{jumpF } f \text{ at-bot} \text{ else } \text{jumpF } f (\text{at-right } (\text{tan } x)))$

$\text{jumpF } (f \circ \text{tan}) (\text{at-left } x) = (\text{if } \cos x = 0$
 $\quad \text{then } \text{jumpF } f \text{ at-top} \text{ else } \text{jumpF } f (\text{at-left } (\text{tan } x)))$

$\langle \text{proof} \rangle$

5.6 Finite jumpFs over an interval

definition *finite-jumpFs*:: $(\text{real} \Rightarrow \text{real}) \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{bool}$ **where**

$\text{finite-jumpFs } f a b = \text{finite } \{x. (\text{jumpF } f (\text{at-left } x) \neq 0 \vee \text{jumpF } f (\text{at-right } x) \neq 0) \wedge a \leq x \wedge x \leq b\}$

lemma *finite-jumpFs-linear-pos*:

assumes $c > 0$

shows $\text{finite-jumpFs } (f \circ (\lambda x. c * x + b)) \text{ lb ub} \longleftrightarrow \text{finite-jumpFs } f (c * \text{lb} + b)$
 $(c * \text{ub} + b)$

$\langle \text{proof} \rangle$

lemma *finite-jumpFs-consts*:

$\text{finite-jumpFs } (\lambda \cdot . c) \text{ lb ub}$

$\langle \text{proof} \rangle$

lemma *finite-jumpFs-combine*:

assumes $\text{finite-jumpFs } f a b$ $\text{finite-jumpFs } f b c$

shows $\text{finite-jumpFs } f a c$

$\langle \text{proof} \rangle$

lemma *finite-jumpFs-subE*:

assumes $\text{finite-jumpFs } f a b$ $a \leq a'$ $b' \leq b$

shows $\text{finite-jumpFs } f a' b'$

<proof>

lemma *finite-Psegments-Re-imp-jumpFs*:

assumes *finite-Psegments* $(\lambda t. \operatorname{Re}(g t - z) = 0)$ *a b continuous-on* $\{a..b\}$ *g*

shows *finite-jumpFs* $(\lambda t. \operatorname{Im}(g t - z)/\operatorname{Re}(g t - z))$ *a b*

<proof>

lemma *finite-ReZ-segments-imp-jumpFs*:

assumes *finite-ReZ-segments* *g z path g*

shows *finite-jumpFs* $(\lambda t. \operatorname{Im}(g t - z)/\operatorname{Re}(g t - z))$ *0 1*

<proof>

5.7 *jumpF* at path ends

definition *jumpF-pathstart*:: $(\operatorname{real} \Rightarrow \operatorname{complex}) \Rightarrow \operatorname{complex} \Rightarrow \operatorname{real}$ **where**

jumpF-pathstart *g z* = *jumpF* $(\lambda t. \operatorname{Im}(g t - z)/\operatorname{Re}(g t - z))$ (*at-right 0*)

definition *jumpF-pathfinish*:: $(\operatorname{real} \Rightarrow \operatorname{complex}) \Rightarrow \operatorname{complex} \Rightarrow \operatorname{real}$ **where**

jumpF-pathfinish *g z* = *jumpF* $(\lambda t. \operatorname{Im}(g t - z)/\operatorname{Re}(g t - z))$ (*at-left 1*)

lemma *jumpF-pathstart-eq-0*:

assumes *path g* $\operatorname{Re}(\operatorname{pathstart} g) \neq \operatorname{Re} z$

shows *jumpF-pathstart* *g z* = 0

<proof>

lemma *jumpF-pathfinish-eq-0*:

assumes *path g* $\operatorname{Re}(\operatorname{pathfinish} g) \neq \operatorname{Re} z$

shows *jumpF-pathfinish* *g z* = 0

<proof>

lemma

shows *jumpF-pathfinish-reversepath*: *jumpF-pathfinish* (*reversepath g*) *z* = *jumpF-pathstart* *g z*

and *jumpF-pathstart-reversepath*: *jumpF-pathstart* (*reversepath g*) *z* = *jumpF-pathfinish* *g z*

<proof>

lemma *jumpF-pathstart-joinpaths[simp]*:

jumpF-pathstart (*g1+++g2*) *z* = *jumpF-pathstart* *g1 z*

<proof>

lemma *jumpF-pathfinish-joinpaths[simp]*:

jumpF-pathfinish (*g1+++g2*) *z* = *jumpF-pathfinish* *g2 z*

<proof>

5.8 Cauchy index

definition *cindex*:: $\operatorname{real} \Rightarrow \operatorname{real} \Rightarrow (\operatorname{real} \Rightarrow \operatorname{real}) \Rightarrow \operatorname{int}$ **where**

cindex *a b f* = $(\sum x \in \{x. \operatorname{jump} f x \neq 0 \wedge a < x \wedge x < b\}. \operatorname{jump} f x)$

definition $cindexE::real \Rightarrow real \Rightarrow (real \Rightarrow real) \Rightarrow real$ **where**
 $cindexE\ a\ b\ f = (\sum x \in \{x. jumpF\ f\ (at-right\ x) \neq 0 \wedge a \leq x \wedge x < b\}. jumpF\ f\ (at-right\ x))$
 $- (\sum x \in \{x. jumpF\ f\ (at-left\ x) \neq 0 \wedge a < x \wedge x \leq b\}. jumpF\ f\ (at-left\ x))$

definition $cindexE-ubd::(real \Rightarrow real) \Rightarrow real$ **where**
 $cindexE-ubd\ f = (\sum x \in \{x. jumpF\ f\ (at-right\ x) \neq 0\}. jumpF\ f\ (at-right\ x))$
 $- (\sum x \in \{x. jumpF\ f\ (at-left\ x) \neq 0\}. jumpF\ f\ (at-left\ x))$

lemma $cindexE-empty$:
 $cindexE\ a\ a\ f = 0$
 $\langle proof \rangle$

lemma $cindex-const$: $cindex\ a\ b\ (\lambda-. c) = 0$
 $\langle proof \rangle$

lemma $cindex-eq-cindex-poly$: $cindex\ a\ b\ (\lambda x. poly\ q\ x / poly\ p\ x) = cindex-poly\ a\ b\ q\ p$
 $\langle proof \rangle$

lemma $cindex-combine$:
assumes $finite:finite\ \{x. jump\ f\ x \neq 0 \wedge a < x \wedge x < c\}$ **and** $a < b < c$
shows $cindex\ a\ c\ f = cindex\ a\ b\ f + jump\ f\ b + cindex\ b\ c\ f$
 $\langle proof \rangle$

lemma $cindexE-combine$:
assumes $finite:finite-jumpFs\ f\ a\ c$ **and** $a \leq b < c$
shows $cindexE\ a\ c\ f = cindexE\ a\ b\ f + cindexE\ b\ c\ f$
 $\langle proof \rangle$

lemma $cindex-linear-comp$:
assumes $c \neq 0$
shows $cindex\ lb\ ub\ (f\ o\ (\lambda x. c*x+b)) = (if\ c > 0$
 $then\ cindex\ (c*lb+b)\ (c*ub+b)\ f$
 $else\ -\ cindex\ (c*ub+b)\ (c*lb+b)\ f)$
 $\langle proof \rangle$

lemma $cindexE-linear-comp$:
assumes $c \neq 0$
shows $cindexE\ lb\ ub\ (f\ o\ (\lambda x. c*x+b)) = (if\ c > 0$
 $then\ cindexE\ (c*lb+b)\ (c*ub+b)\ f$
 $else\ -\ cindexE\ (c*ub+b)\ (c*lb+b)\ f)$
 $\langle proof \rangle$

lemma $cindexE-cong$:
assumes $finite\ s$ **and** $fg-eq:\bigwedge x. \llbracket a < x; x < b; x \notin s \rrbracket \implies f\ x = g\ x$
shows $cindexE\ a\ b\ f = cindexE\ a\ b\ g$

<proof>

lemma *cindexE-constI*:

assumes $\bigwedge t. \llbracket a < t; t < b \rrbracket \implies f t = c$

shows $cindexE\ a\ b\ f = 0$

<proof>

lemma *cindex-eq-cindexE-divide*:

fixes $f\ g::real \Rightarrow real$

defines $h \equiv (\lambda x. f\ x/g\ x)$

assumes $a < b$ **and**

finite-fg: $finite\ \{x. (f\ x=0 \vee g\ x=0) \wedge a \leq x \wedge x \leq b\}$ **and**

g-imp-f: $\forall x \in \{a..b\}. g\ x=0 \longrightarrow f\ x \neq 0$ **and**

f-cont: *continuous-on* $\{a..b\}$ f **and**

g-cont: *continuous-on* $\{a..b\}$ g

shows $cindexE\ a\ b\ h = jumpF\ h\ (at-right\ a) + cindex\ a\ b\ h - jumpF\ h\ (at-left\ b)$

<proof>

5.9 Cauchy index along a path

definition *cindex-path*:: $(real \Rightarrow complex) \Rightarrow complex \Rightarrow int$ **where**

$cindex-path\ g\ z = cindex\ 0\ 1\ (\lambda t. Im\ (g\ t - z) / Re\ (g\ t - z))$

definition *cindex-pathE*:: $(real \Rightarrow complex) \Rightarrow complex \Rightarrow real$ **where**

$cindex-pathE\ g\ z = cindexE\ 0\ 1\ (\lambda t. Im\ (g\ t - z) / Re\ (g\ t - z))$

lemma *cindex-pathE-point*: $cindex-pathE\ (linepath\ a\ a)\ b = 0$

<proof>

lemma *cindex-path-reversepath*:

$cindex-path\ (reversepath\ g)\ z = -\ cindex-path\ g\ z$

<proof>

lemma *cindex-pathE-reversepath*: $cindex-pathE\ (reversepath\ g)\ z = -\ cindex-pathE\ g\ z$

<proof>

lemma *cindex-pathE-reversepath'*: $cindex-pathE\ g\ z = -\ cindex-pathE\ (reversepath\ g)\ z$

<proof>

lemma *cindex-pathE-joinpaths*:

assumes $g1$: *finite-ReZ-segments* $g1\ z$ **and** $g2$: *finite-ReZ-segments* $g2\ z$ **and**

$path\ g1\ path\ g2\ pathfinish\ g1 = pathstart\ g2$

shows $cindex-pathE\ (g1+++g2)\ z = cindex-pathE\ g1\ z + cindex-pathE\ g2\ z$

<proof>

lemma *cindex-pathE-constI*:

assumes $\bigwedge t. \llbracket 0 < t; t < 1 \rrbracket \implies g\ t = c$
shows $\text{cindex-pathE } g\ z = 0$
 $\langle \text{proof} \rangle$

lemma *cindex-pathE-subpath-combine*:

assumes $g::\text{finite-ReZ-segments } g$ **and** $\text{path } g$ **and**

$0 \leq a \leq b \leq c \leq 1$

shows $\text{cindex-pathE } (\text{subpath } a\ b\ g)\ z + \text{cindex-pathE } (\text{subpath } b\ c\ g)\ z$
 $= \text{cindex-pathE } (\text{subpath } a\ c\ g)\ z$

$\langle \text{proof} \rangle$

lemma *cindex-pathE-shiftpath*:

assumes $\text{finite-ReZ-segments } g\ z$ $s \in \{0..1\}$ $\text{path } g$ **and** $\text{loop:pathfinish } g = \text{pathstart } g$

shows $\text{cindex-pathE } (\text{shiftpath } s\ g)\ z = \text{cindex-pathE } g\ z$

$\langle \text{proof} \rangle$

5.10 Cauchy's Index Theorem

theorem *winding-number-cindex-pathE-aux*:

fixes $g::\text{real} \Rightarrow \text{complex}$

assumes $\text{finite-ReZ-segments } g\ z$ **and** $\text{valid-path } g\ z \notin \text{path-image } g$ **and**

$\text{Re-ends:Re } (g\ 1) = \text{Re } z$ $\text{Re } (g\ 0) = \text{Re } z$

shows $2 * \text{Re}(\text{winding-number } g\ z) = - \text{cindex-pathE } g\ z$

$\langle \text{proof} \rangle$

theorem *winding-number-cindex-pathE*:

fixes $g::\text{real} \Rightarrow \text{complex}$

assumes $\text{finite-ReZ-segments } g\ z$ **and** $\text{valid-path } g\ z \notin \text{path-image } g$ **and**

$\text{loop: pathfinish } g = \text{pathstart } g$

shows $\text{winding-number } g\ z = - \text{cindex-pathE } g\ z / 2$

$\langle \text{proof} \rangle$

REMARK: The usual statement of Cauchy's Index theorem (i.e. Analytic Theory of Polynomials (2002): Theorem 11.1.3) is about the equality between the number of polynomial roots and the Cauchy index, which is the joint application of $\llbracket \text{finite-ReZ-segments } ?g\ ?z; \text{valid-path } ?g; ?z \notin \text{path-image } ?g; \text{pathfinish } ?g = \text{pathstart } ?g \rrbracket \implies \text{winding-number } ?g\ ?z = \text{complex-of-real } (- \text{cindex-pathE } ?g\ ?z / 2)$ and $\llbracket \text{open } ?s; \text{connected } ?s; ?f \text{ holomorphic-on } ?s - ?poles; ?h \text{ holomorphic-on } ?s; \text{valid-path } ?g; \text{pathfinish } ?g = \text{pathstart } ?g; \text{path-image } ?g \subseteq ?s - \{w. ?f\ w = 0 \vee w \in ?poles\}; \forall z. z \notin ?s \longrightarrow \text{winding-number } ?g\ z = 0; \text{finite } \{w. ?f\ w = 0 \vee w \in ?poles\}; \forall p \in ?poles. \text{is-pole } ?f\ p \rrbracket \implies \text{contour-integral } ?g\ (\lambda x. \text{deriv } ?f\ x * ?h\ x / ?f\ x) = \text{complex-of-real } (2 * \text{pi}) * i * (\sum p \in \{w. ?f\ w = 0 \vee w \in ?poles\}. \text{winding-number } ?g\ p * ?h\ p * \text{complex-of-int } (\text{zorder } ?f\ p))$.

end

6 Evaluate winding numbers by calculating Cauchy indices

```

theory Winding-Number-Eval imports
  Cauchy-Index-Theorem
  HOL-Eisbach.Eisbach-Tools
begin

```

6.1 Misc

```

lemma not-on-closed-segmentI:
  fixes z::'a::euclidean-space
  assumes  $\text{norm } (z - a) *_{\mathbb{R}} (b - z) \neq \text{norm } (b - z) *_{\mathbb{R}} (z - a)$ 
  shows  $z \notin \text{closed-segment } a \ b$ 
   $\langle \text{proof} \rangle$ 

```

```

lemma not-on-closed-segmentI-complex:
  fixes z::complex
  assumes  $(\text{Re } b - \text{Re } z) * (\text{Im } z - \text{Im } a) \neq (\text{Im } b - \text{Im } z) * (\text{Re } z - \text{Re } a)$ 
  shows  $z \notin \text{closed-segment } a \ b$ 
   $\langle \text{proof} \rangle$ 

```

6.2 finite intersection with the two axes

```

definition finite-axes-cross:: $(\text{real} \Rightarrow \text{complex}) \Rightarrow \text{complex} \Rightarrow \text{bool}$  where
  finite-axes-cross g z = finite  $\{t. (\text{Re } (g \ t - z) = 0 \vee \text{Im } (g \ t - z) = 0) \wedge 0 \leq t \wedge t \leq 1\}$ 

```

```

lemma finite-cross-intros:
   $\llbracket \text{Re } a \neq \text{Re } z \vee \text{Re } b \neq \text{Re } z; \text{Im } a \neq \text{Im } z \vee \text{Im } b \neq \text{Im } z \rrbracket \Longrightarrow \text{finite-axes-cross } (\text{linepath } a \ b) \ z$ 
   $\llbracket st \neq tt; r \neq 0 \rrbracket \Longrightarrow \text{finite-axes-cross } (\text{part-circlepath } z0 \ r \ st \ tt) \ z$ 
   $\llbracket \text{finite-axes-cross } g1 \ z; \text{finite-axes-cross } g2 \ z \rrbracket \Longrightarrow \text{finite-axes-cross } (g1+++g2) \ z$ 
   $\langle \text{proof} \rangle$ 

```

```

lemma cindex-path-joinpaths:
  assumes finite-axes-cross g1 z finite-axes-cross g2 z
  and path g1 path g2 pathfinish g1 = pathstart g2 pathfinish g1  $\neq z$ 
  shows  $\text{cindex-path } (g1+++g2) \ z = \text{cindex-path } g1 \ z + \text{jumpF-pathstart } g2 \ z$ 
   $\quad - \text{jumpF-pathfinish } g1 \ z + \text{cindex-path } g2 \ z$ 
   $\langle \text{proof} \rangle$ 

```

6.3 More lemmas related *cindex-pathE* / *jumpF-pathstart* / *jumpF-pathfinish*

```

lemma cindex-pathE-linepath:
  assumes  $z \notin \text{closed-segment } a \ b$ 
  shows  $\text{cindex-pathE } (\text{linepath } a \ b) \ z = ($ 
     $\text{let } c1 = \text{Re } a - \text{Re } z;$ 
     $\quad c2 = \text{Re } b - \text{Re } z;$ 

```

$c3 = \text{Im } a * \text{Re } b + \text{Re } z * \text{Im } b + \text{Im } z * \text{Re } a - \text{Im } z * \text{Re } b - \text{Im } b * \text{Re } a - \text{Re } z * \text{Im } a;$
 $d1 = \text{Im } a - \text{Im } z;$
 $d2 = \text{Im } b - \text{Im } z$
in if $(c1 > 0 \wedge c2 < 0) \vee (c1 < 0 \wedge c2 > 0)$ then
 (if $c3 > 0$ then 1 else -1)
else
 (if $(c1 = 0 \leftrightarrow c2 \neq 0) \wedge (c1 = 0 \rightarrow d1 \neq 0) \wedge (c2 = 0 \rightarrow d2 \neq 0)$ then
 if $(c1 = 0 \wedge (c2 > 0 \leftrightarrow d1 > 0)) \vee (c2 = 0 \wedge (c1 > 0 \leftrightarrow d2 < 0))$ then
1/2 else -1/2
 else 0))
⟨proof⟩

lemma *cindex-path-linepath:*

assumes $z \notin \text{path-image } (\text{linepath } a \ b)$
shows $\text{cindex-path } (\text{linepath } a \ b) \ z = (
 \text{let } c1 = \text{Re}(a) - \text{Re}(z) ; c2 = \text{Re}(b) - \text{Re}(z) ;
 c3 = \text{Im}(a) * \text{Re}(b) + \text{Re}(z) * \text{Im}(b) + \text{Im}(z) * \text{Re}(a) - \text{Im}(z) * \text{Re}(b) - \text{Im}(b) * \text{Re}(a)
- \text{Re}(z) * \text{Im}(a)
in if $(c1 > 0 \wedge c2 < 0) \vee (c1 < 0 \wedge c2 > 0)$ then (if $c3 > 0$ then 1 else -1) else 0)$
⟨proof⟩

lemma *cindex-pathE-part-circlepath:*

assumes $\text{cmod } (z - z0) \neq r$ **and** $r > 0 \ 0 \leq st \ st < tt \ tt \leq 2 * \pi$
shows $\text{cindex-pathE } (\text{part-circlepath } z \ r \ st \ tt) \ z0 = (
 \text{if } |\text{Re } z - \text{Re } z0| < r \text{ then}$
 (let
 $\vartheta = \arccos ((\text{Re } z0 - \text{Re } z) / r);$
 $\beta = 2 * \pi - \vartheta$
in
 $\text{jumpF-pathstart } (\text{part-circlepath } z \ r \ st \ tt) \ z0$
 +
 (if $st < \vartheta \wedge \vartheta < tt$ then if $r * \sin \vartheta + \text{Im } z > \text{Im } z0$ then -1 else 1 else 0)
 +
 (if $st < \beta \wedge \beta < tt$ then if $r * \sin \beta + \text{Im } z > \text{Im } z0$ then 1 else -1 else 0)
 -
 $\text{jumpF-pathfinish } (\text{part-circlepath } z \ r \ st \ tt) \ z0$
)
 else
 if $|\text{Re } z - \text{Re } z0| = r$ then
 $\text{jumpF-pathstart } (\text{part-circlepath } z \ r \ st \ tt) \ z0$
 - $\text{jumpF-pathfinish } (\text{part-circlepath } z \ r \ st \ tt) \ z0$
 else 0
)
⟨proof⟩

lemma *jumpF-pathstart-part-circlepath:*

assumes $st < tt \ r > 0 \ \text{cmod } (z - z0) \neq r$

shows $\text{jumpF-pathstart } (\text{part-circlepath } z \ r \ st \ tt) \ z0 = ($
 if $r * \cos st + \text{Re } z - \text{Re } z0 = 0$ then
 (let
 $\Delta = r * \sin st + \text{Im } z - \text{Im } z0$
 in
 if $(\sin st > 0 \vee \cos st = 1) \wedge \Delta < 0$
 $\vee (\sin st < 0 \vee \cos st = -1) \wedge \Delta > 0$ then
 $1/2$
 else
 $-1/2$)
 else 0)
 ⟨proof⟩

lemma $\text{jumpF-pathfinish-part-circlepath}$:
assumes $st < tt \ r > 0 \ cmod (z - z0) \neq r$
shows $\text{jumpF-pathfinish } (\text{part-circlepath } z \ r \ st \ tt) \ z0 = ($
 if $r * \cos tt + \text{Re } z - \text{Re } z0 = 0$ then
 (let
 $\Delta = r * \sin tt + \text{Im } z - \text{Im } z0$
 in
 if $(\sin tt > 0 \vee \cos tt = -1) \wedge \Delta < 0$
 $\vee (\sin tt < 0 \vee \cos tt = 1) \wedge \Delta > 0$ then
 $-1/2$
 else
 $1/2$)
 else 0)
 ⟨proof⟩

lemma
fixes $z0 \ z :: \text{complex}$ **and** $r :: \text{real}$
defines $\text{upper} \equiv \text{cindex-pathE } (\text{part-circlepath } z \ r \ 0 \ pi) \ z0$
and $\text{lower} \equiv \text{cindex-pathE } (\text{part-circlepath } z \ r \ pi \ (2 * pi)) \ z0$
shows $\text{cindex-pathE-circlepath-upper}$:
 $\llbracket cmod (z0 - z) < r \rrbracket \implies \text{upper} = -1$
 $\llbracket \text{Im } (z0 - z) > r; |\text{Re } (z0 - z)| < r \rrbracket \implies \text{upper} = 1$
 $\llbracket \text{Im } (z0 - z) < -r; |\text{Re } (z0 - z)| < r \rrbracket \implies \text{upper} = -1$
 $\llbracket \text{Re } (z0 - z) > r; r > 0 \rrbracket \implies \text{upper} = 0$
and $\text{cindex-pathE-circlepath-lower}$:
 $\llbracket cmod (z0 - z) < r \rrbracket \implies \text{lower} = -1$
 $\llbracket \text{Im } (z0 - z) > r; |\text{Re } (z0 - z)| < r \rrbracket \implies \text{lower} = -1$
 $\llbracket \text{Im } (z0 - z) < -r; |\text{Re } (z0 - z)| < r \rrbracket \implies \text{lower} = 1$
 $\llbracket \text{Re } (z0 - z) > r; r > 0 \rrbracket \implies \text{lower} = 0$
 ⟨proof⟩

lemma $\text{jumpF-pathstart-linepath}$:
 $\text{jumpF-pathstart } (\text{linepath } a \ b) \ z =$
 (if $\text{Re } a = \text{Re } z \wedge \text{Im } a \neq \text{Im } z \wedge \text{Re } b \neq \text{Re } a$ then
 if $(\text{Im } a > \text{Im } z \wedge \text{Re } b > \text{Re } a) \vee (\text{Im } a < \text{Im } z \wedge \text{Re } b < \text{Re } a)$ then $1/2$ else
 $-1/2$)

else 0)
 ⟨proof⟩

lemma *jumpF-pathfinish-linepath*:

jumpF-pathfinish (linepath a b) z =
 (if Re b = Re z ∧ Im b ≠ Im z ∧ Re b ≠ Re a then
 if (Im b > Im z ∧ Re a > Re b) ∨ (Im b < Im z ∧ Re a < Re b) then 1/2 else
 -1/2
 else 0)
 ⟨proof⟩

6.4 Setting up the method for evaluating winding numbers

lemma *pathfinish-pathstart-partcirclepath-simps*:

pathstart (part-circlepath z0 r (3*pi/2) tt) = z0 - Complex 0 r
pathstart (part-circlepath z0 r (2*pi) tt) = z0 + r
pathfinish (part-circlepath z0 r st (3*pi/2)) = z0 - Complex 0 r
pathfinish (part-circlepath z0 r st (2*pi)) = z0 + r
pathstart (part-circlepath z0 r 0 tt) = z0 + r
pathstart (part-circlepath z0 r (pi/2) tt) = z0 + Complex 0 r
pathstart (part-circlepath z0 r (pi) tt) = z0 - r
pathfinish (part-circlepath z0 r st 0) = z0 + r
pathfinish (part-circlepath z0 r st (pi/2)) = z0 + Complex 0 r
pathfinish (part-circlepath z0 r st (pi)) = z0 - r
 ⟨proof⟩

lemma *winding-eq-intro*:

finite-ReZ-segments g z ⇒
valid-path g ⇒
 z ∉ *path-image* g ⇒
pathfinish g = *pathstart* g ⇒
 - of-real(*cindex-pathE* g z) = 2*n ⇒
winding-number g z = (n::complex)
 ⟨proof⟩

named-theorems *winding-intros* and *winding-simps*

lemmas [*winding-intros*] =
finite-ReZ-segments-joinpaths
valid-path-join
path-join-imp
not-in-path-image-join

lemmas [*winding-simps*] =
finite-ReZ-segments-linepath
finite-ReZ-segments-part-circlepath
jumpF-pathfinish-joinpaths
jumpF-pathstart-joinpaths
pathfinish-linepath

```

pathstart-linepath
pathfinish-join
pathstart-join
valid-path-linepath
valid-path-part-circlepath
path-part-circlepath
Re-complex-of-real
Im-complex-of-real
of-real-linepath
pathfinish-pathstart-partcirclepath-simps

```

```

method rep-subst =
  (subst cindex-pathE-joinpaths; rep-subst)?

```

The method "eval_winding" $1::'a$ will try to simplify of the form *winding-number* $g z = n$ where n is an integer and g is a closed path comprised of *linepath*, *part-circlepath* and (+++).

Suppose $g = l1 +++ l2$, usually, the key behind the success of this framework is whether we can prove $z \notin \text{path-image } l1$, $z \notin \text{path-image } l2$ and calculate *cindex-pathE* $l1 z$ and *cindex-pathE* $l2 z$.

```

method eval-winding =
  ((rule-tac winding-eq-intro;
    rep-subst
  )
  , auto simp only:winding-simps del:notI intro!:winding-intros
  , tactic <distinct-subgoals-tac>)

```

end

7 Some examples of applying the method winding_eval

```

theory Winding-Number-Eval-Examples imports Winding-Number-Eval
begin

```

```

lemma example1:
  assumes  $R > 1$ 
  shows winding-number (part-circlepath 0 R 0 pi +++ linepath (-R) R) i = 1
  <proof>

```

```

lemma example2:
  assumes  $R > 1$ 
  shows winding-number (part-circlepath 0 R 0 pi +++ linepath (-R) R) (-i) =
  0
  <proof>

```

```

lemma example3:
  fixes lb ub z :: complex

```

```

defines rec  $\equiv$  linepath lb (Complex (Re ub) (Im lb)) +++ linepath (Complex
(Re ub) (Im lb)) ub
      +++ linepath ub (Complex (Re lb) (Im ub)) +++ linepath (Complex
(Re lb) (Im ub)) lb
assumes order-asms:Re lb < Re z Re z < Re ub Im lb < Im z Im z < Im ub
shows winding-number rec z = 1
<proof>

```

end

8 Acknowledgements

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References

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