

Evaluate Winding Numbers through Cauchy Indices

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Abstract

In complex analysis, the winding number measures the number of times a path (counterclockwise) winds around a point, while the Cauchy index can approximate how the path winds. This entry provides a formalisation of the Cauchy index, which is then shown to be related to the winding number. In addition, this entry also offers a tactic that enables users to evaluate the winding number by calculating Cauchy indices. The connection between the winding number and the Cauchy index can be found in the literature [1] [2, Chapter 11].

1 Some useful lemmas in topology

theory *Missing-Topology* **imports** *HOL-Analysis.Multivariate-Analysis*
begin

1.1 Misc

lemma *open-times-image*:

fixes $S::'a::\text{real-normed-field}$ *set*

assumes *open S c* $\neq 0$

shows *open* $((\text{c} \cdot S) \cdot S)$

<proof>

lemma *image-linear-greaterThan*:

fixes $x::'a::\text{linordered-field}$

assumes $c \neq 0$

shows $((\lambda x. c \cdot x + b) \cdot \{x < ..\}) = (\text{if } c > 0 \text{ then } \{c \cdot x + b < ..\} \text{ else } \{.. < c \cdot x + b\})$

<proof>

lemma *image-linear-lessThan*:

fixes $x::'a::\text{linordered-field}$

assumes $c \neq 0$

shows $((\lambda x. c \cdot x + b) \cdot \{.. < x\}) = (\text{if } c > 0 \text{ then } \{.. < c \cdot x + b\} \text{ else } \{c \cdot x + b < ..\})$

<proof>

lemma *continuous-on-neq-split*:

fixes $f :: 'a::\text{linear-continuum-topology} \Rightarrow 'b::\text{linorder-topology}$

assumes $\forall x \in s. f x \neq y$ *continuous-on s f connected s*
shows $(\forall x \in s. f x > y) \vee (\forall x \in s. f x < y)$
 ⟨proof⟩

lemma
fixes $f::'a::\text{linorder-topology} \Rightarrow 'b::\text{topological-space}$
assumes *continuous-on {a..b} f a < b*
shows *continuous-on-at-left:continuous (at-left b) f*
and *continuous-on-at-right:continuous (at-right a) f*
 ⟨proof⟩

1.2 More about *eventually*

lemma *eventually-comp-filtermap*:
eventually (P o f) F \longleftrightarrow eventually P (filtermap f F)
 ⟨proof⟩

lemma *eventually-uminus-at-top-at-bot*:
fixes $P::'a::\{\text{ordered-ab-group-add,linorder}\} \Rightarrow \text{bool}$
shows *eventually (P o uminus) at-bot \longleftrightarrow eventually P at-top*
eventually (P o uminus) at-top \longleftrightarrow eventually P at-bot
 ⟨proof⟩

lemma *eventually-at-infinityI*:
fixes $P::'a::\text{real-normed-vector} \Rightarrow \text{bool}$
assumes $\bigwedge x. c \leq \text{norm } x \implies P x$
shows *eventually P at-infinity*
 ⟨proof⟩

lemma *eventually-at-bot-linorderI*:
fixes $c::'a::\text{linorder}$
assumes $\bigwedge x. x \leq c \implies P x$
shows *eventually P at-bot*
 ⟨proof⟩

lemma *eventually-times-inverse-1*:
fixes $f::'a \Rightarrow 'b::\{\text{field,t2-space}\}$
assumes $(f \longrightarrow c) F c \neq 0$
shows $\forall_F x \text{ in } F. \text{inverse } (f x) * f x = 1$
 ⟨proof⟩

1.3 More about *filtermap*

lemma *filtermap-linear-at-within*:
assumes *bij f and cont: isCont f a and open-map: $\bigwedge S. \text{open } S \implies \text{open } (f'S)$*
shows *filtermap f (at a within S) = at (f a) within f'S*
 ⟨proof⟩

lemma *filtermap-at-bot-linear-eq*:
fixes $c::'a::\text{linordered-field}$

assumes $c \neq 0$
shows $\text{filtermap } (\lambda x. x * c + b) \text{ at-bot} = (\text{if } c > 0 \text{ then at-bot else at-top})$
 $\langle \text{proof} \rangle$

lemma *filtermap-linear-at-left*:
fixes $c :: 'a :: \{\text{linordered-field, linorder-topology, real-normed-field}\}$
assumes $c \neq 0$
shows $\text{filtermap } (\lambda x. c * x + b) (\text{at-left } x) = (\text{if } c > 0 \text{ then at-left } (c * x + b) \text{ else at-right } (c * x + b))$
 $\langle \text{proof} \rangle$

lemma *filtermap-linear-at-right*:
fixes $c :: 'a :: \{\text{linordered-field, linorder-topology, real-normed-field}\}$
assumes $c \neq 0$
shows $\text{filtermap } (\lambda x. c * x + b) (\text{at-right } x) = (\text{if } c > 0 \text{ then at-right } (c * x + b) \text{ else at-left } (c * x + b))$
 $\langle \text{proof} \rangle$

lemma *filtermap-at-top-linear-eq*:
fixes $c :: 'a :: \text{linordered-field}$
assumes $c \neq 0$
shows $\text{filtermap } (\lambda x. x * c + b) \text{ at-top} = (\text{if } c > 0 \text{ then at-top else at-bot})$
 $\langle \text{proof} \rangle$

lemma *filtermap-nhds-open-map*:
assumes $\text{cont: isCont } f \ a$
and $\text{open-map: } \bigwedge S. \text{open } S \implies \text{open } (f'S)$
shows $\text{filtermap } f \ (\text{nhds } a) = \text{nhds } (f \ a)$
 $\langle \text{proof} \rangle$

1.4 More about *filterlim*

lemma *filterlim-at-infinity-times*:
fixes $f :: 'a \Rightarrow 'b :: \text{real-normed-field}$
assumes $\text{filterlim } f \ \text{at-infinity } F \ \text{filterlim } g \ \text{at-infinity } F$
shows $\text{filterlim } (\lambda x. f \ x * g \ x) \ \text{at-infinity } F$
 $\langle \text{proof} \rangle$

lemma *filterlim-at-top-at-bot[elim]*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{unbounded-dense-linorder and } F :: 'a \text{ filter}\}$
assumes $\text{top: filterlim } f \ \text{at-top } F$ **and** $\text{bot: filterlim } f \ \text{at-bot } F$ **and** $F \neq \text{bot}$
shows False
 $\langle \text{proof} \rangle$

lemma *filterlim-at-top-nhds[elim]*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{unbounded-dense-linorder, order-topology}\}$ **and** $F :: 'a \text{ filter}$
assumes $\text{top: filterlim } f \ \text{at-top } F$ **and** $\text{tendsto: } (f \longrightarrow c) \ F$ **and** $F \neq \text{bot}$
shows False

<proof>

lemma *filterlim-at-bot-nhds*[*elim*]:

fixes $f::'a \Rightarrow 'b::\{\text{unbounded-dense-linorder}, \text{order-topology}\}$ **and** $F::'a$ *filter*

assumes *top:filterlim f at-bot F* **and** *tendsto: (f \longrightarrow c) F* **and** $F \neq \text{bot}$

shows *False*

<proof>

lemma *filterlim-at-top-linear-iff*:

fixes $f::'a::\text{linordered-field} \Rightarrow 'b$

assumes $c \neq 0$

shows $(\text{LIM } x \text{ at-top. } f (x * c + b) :> F2) \iff (\text{if } c > 0 \text{ then } (\text{LIM } x \text{ at-top. } f x :> F2))$

$\text{else } (\text{LIM } x \text{ at-bot. } f x :> F2))$

<proof>

lemma *filterlim-at-bot-linear-iff*:

fixes $f::'a::\text{linordered-field} \Rightarrow 'b$

assumes $c \neq 0$

shows $(\text{LIM } x \text{ at-bot. } f (x * c + b) :> F2) \iff (\text{if } c > 0 \text{ then } (\text{LIM } x \text{ at-bot. } f x :> F2))$

$\text{else } (\text{LIM } x \text{ at-top. } f x :> F2))$

<proof>

lemma *filterlim-tendsto-add-at-top-iff*:

assumes $f: (f \longrightarrow c) F$

shows $(\text{LIM } x F. (f x + g x :: \text{real}) :> \text{at-top}) \iff (\text{LIM } x F. g x :> \text{at-top})$

<proof>

lemma *filterlim-tendsto-add-at-bot-iff*:

fixes $c::\text{real}$

assumes $f: (f \longrightarrow c) F$

shows $(\text{LIM } x F. f x + g x :> \text{at-bot}) \iff (\text{LIM } x F. g x :> \text{at-bot})$

<proof>

lemma *tendsto-inverse-0-at-infinity*:

$\text{LIM } x F. f x :> \text{at-infinity} \implies ((\lambda x. \text{inverse } (f x) :: \text{real}) \longrightarrow 0) F$

<proof>

lemma *filterlim-at-infinity-divide-iff*:

fixes $f::'a \Rightarrow 'b::\text{real-normed-field}$

assumes $(f \longrightarrow c) F$ $c \neq 0$

shows $(\text{LIM } x F. f x / g x :> \text{at-infinity}) \iff (\text{LIM } x F. g x :> \text{at } 0)$

<proof>

lemma *filterlim-tendsto-pos-mult-at-top-iff*:

fixes $f::'a \Rightarrow \text{real}$

assumes $(f \longrightarrow c) F$ **and** $0 < c$
shows $(LIM\ x\ F.\ (f\ x\ * \ g\ x) :> \text{at-top}) \longleftrightarrow (LIM\ x\ F.\ g\ x :> \text{at-top})$
 $\langle \text{proof} \rangle$

lemma *filterlim-tendsto-pos-mult-at-bot-iff*:
fixes $c :: \text{real}$
assumes $(f \longrightarrow c) F$ $0 < c$
shows $(LIM\ x\ F.\ f\ x\ * \ g\ x :> \text{at-bot}) \longleftrightarrow \text{filterlim}\ g\ \text{at-bot}\ F$
 $\langle \text{proof} \rangle$

lemma *filterlim-tendsto-neg-mult-at-top-iff*:
fixes $f :: 'a \Rightarrow \text{real}$
assumes $(f \longrightarrow c) F$ **and** $c < 0$
shows $(LIM\ x\ F.\ (f\ x\ * \ g\ x) :> \text{at-top}) \longleftrightarrow (LIM\ x\ F.\ g\ x :> \text{at-bot})$
 $\langle \text{proof} \rangle$

lemma *filterlim-tendsto-neg-mult-at-bot-iff*:
fixes $c :: \text{real}$
assumes $(f \longrightarrow c) F$ $0 > c$
shows $(LIM\ x\ F.\ f\ x\ * \ g\ x :> \text{at-bot}) \longleftrightarrow \text{filterlim}\ g\ \text{at-top}\ F$
 $\langle \text{proof} \rangle$

lemma *Lim-add*:
fixes $f\ g :: 'a :: \{t2\text{-space}, \text{topological-monoid-add}\}$
assumes $\exists y.\ (f \longrightarrow y) F$ **and** $\exists y.\ (g \longrightarrow y) F$ **and** $F \neq \text{bot}$
shows $\text{Lim}\ F\ f\ +\ \text{Lim}\ F\ g\ =\ \text{Lim}\ F\ (\lambda x.\ f\ x\ +\ g\ x)$
 $\langle \text{proof} \rangle$

1.5 Isolate and discrete

definition (**in** *topological-space*) *isolate*:: $'a \text{ set} \Rightarrow \text{bool}$ (**infixr** *isolate* 60)
where $x\ \text{isolate}\ S \longleftrightarrow (x \in S \wedge (\exists T.\ \text{open}\ T \wedge T \cap S = \{x\}))$

definition (**in** *topological-space*) *discrete*:: $'a \text{ set} \Rightarrow \text{bool}$
where $\text{discrete}\ S \longleftrightarrow (\forall x \in S.\ x\ \text{isolate}\ S)$

definition (**in** *metric-space*) *uniform-discrete* :: $'a \text{ set} \Rightarrow \text{bool}$ **where**
uniform-discrete $S \longleftrightarrow (\exists e > 0.\ \forall x \in S.\ \forall y \in S.\ \text{dist}\ x\ y < e \longrightarrow x = y)$

lemma *uniformI1*:
assumes $e > 0 \wedge x\ y.\ \llbracket x \in S; y \in S; \text{dist}\ x\ y < e \rrbracket \Longrightarrow x = y$
shows *uniform-discrete* S
 $\langle \text{proof} \rangle$

lemma *uniformI2*:
assumes $e > 0 \wedge x\ y.\ \llbracket x \in S; y \in S; x \neq y \rrbracket \Longrightarrow \text{dist}\ x\ y \geq e$
shows *uniform-discrete* S
 $\langle \text{proof} \rangle$

lemma *isolate-islimpt-iff*: $(x \text{ isolate } S) \longleftrightarrow (\neg (x \text{ islimpt } S) \wedge x \in S)$
<proof>

lemma *isolate-dist-Ex-iff*:
fixes $x :: 'a :: \text{metric-space}$
shows $x \text{ isolate } S \longleftrightarrow (x \in S \wedge (\exists e > 0. \forall y \in S. \text{dist } x \ y < e \longrightarrow y = x))$
<proof>

lemma *discrete-empty[simp]*: *discrete* {}
<proof>

lemma *uniform-discrete-empty[simp]*: *uniform-discrete* {}
<proof>

lemma *isolate-insert*:
fixes $x :: 'a :: \text{t1-space}$
shows $x \text{ isolate } (\text{insert } a \ S) \longleftrightarrow x \text{ isolate } S \vee (x = a \wedge \neg (x \text{ islimpt } S))$
<proof>

lemma *uniform-discrete-imp-closed*:
uniform-discrete $S \implies \text{closed } S$
<proof>

lemma *uniform-discrete-imp-discrete*:
uniform-discrete $S \implies \text{discrete } S$
<proof>

lemma *isolate-subset*: $x \text{ isolate } S \implies T \subseteq S \implies x \in T \implies x \text{ isolate } T$
<proof>

lemma *discrete-subset[elim]*: *discrete* $S \implies T \subseteq S \implies \text{discrete } T$
<proof>

lemma *uniform-discrete-subset[elim]*: *uniform-discrete* $S \implies T \subseteq S \implies \text{uniform-discrete } T$
<proof>

lemma *continuous-on-discrete*: *discrete* $S \implies \text{continuous-on } S \ f$
<proof>

lemma *uniform-discrete-insert*:
fixes $S :: 'a :: \text{euclidean-space set}$
shows $\text{uniform-discrete } (\text{insert } a \ S) \longleftrightarrow \text{uniform-discrete } S$
<proof>

lemma *discrete-compact-finite-iff*:

fixes $S :: 'a::t1\text{-space set}$
shows $\text{discrete } S \wedge \text{compact } S \longleftrightarrow \text{finite } S$
 $\langle \text{proof} \rangle$

lemma *uniform-discrete-finite-iff*:
fixes $S :: 'a::\text{heine-borel set}$
shows $\text{uniform-discrete } S \wedge \text{bounded } S \longleftrightarrow \text{finite } S$
 $\langle \text{proof} \rangle$

lemma *uniform-discrete-image-scale*:
fixes $f :: 'a::\text{euclidean-space} \Rightarrow 'b::\text{euclidean-space}$
assumes $\text{uniform-discrete } S$ **and** $\text{dist}:\forall x\in S. \forall y\in S. \text{dist } x \ y = c * \text{dist } (f \ x) \ (f \ y)$
shows $\text{uniform-discrete } (f \ ' S)$
 $\langle \text{proof} \rangle$

end

2 Some useful lemmas in algebra

theory *Missing-Algebraic imports*
HOL-Computational-Algebra.Polynomial-Factorial
HOL-Computational-Algebra.Fundamental-Theorem-Algebra
HOL-Complex-Analysis.Complex-Analysis
Missing-Topology
Budan-Fourier.BF-Misc
begin

2.1 Misc

lemma *poly-holomorphic-on[simp]*:
 $(\text{poly } p) \text{ holomorphic-on } s$
 $\langle \text{proof} \rangle$

lemma *order-zorder*:
fixes $p::\text{complex poly}$ **and** $z::\text{complex}$
assumes $p \neq 0$
shows $\text{order } z \ p = \text{nat } (\text{zorder } (\text{poly } p) \ z)$
 $\langle \text{proof} \rangle$

lemma *pcompose-pCons-0*: $\text{pcompose } p \ [:a:] = [:\text{poly } p \ a:]$
 $\langle \text{proof} \rangle$

lemma *pcompose-coeff-0*:
 $\text{coeff } (\text{pcompose } p \ q) \ 0 = \text{poly } p \ (\text{coeff } q \ 0)$
 $\langle \text{proof} \rangle$

lemma *poly-field-differentiable-at*[simp]:
poly p field-differentiable (at x within s)
 ⟨proof⟩

lemma *deriv-pderiv*:
deriv (poly p) = poly (pderiv p)
 ⟨proof⟩

lemma *lead-coeff-map-poly-nz*:
assumes *f (lead-coeff p) ≠ 0 f 0 = 0*
shows *lead-coeff (map-poly f p) = f (lead-coeff p)*
 ⟨proof⟩

lemma *filterlim-poly-at-infinity*:
fixes *p::'a::real-normed-field poly*
assumes *degree p > 0*
shows *filterlim (poly p) at-infinity at-infinity*
 ⟨proof⟩

lemma *poly-divide-tendsto-aux*:
fixes *p::'a::real-normed-field poly*
shows *((λx. poly p x / x^(degree p)) → lead-coeff p) at-infinity*
 ⟨proof⟩

lemma *filterlim-power-at-infinity*:
assumes *n ≠ 0*
shows *filterlim (λx::'a::real-normed-field. xⁿ) at-infinity at-infinity*
 ⟨proof⟩

lemma *poly-divide-tendsto-0-at-infinity*:
fixes *p::'a::real-normed-field poly*
assumes *degree p > degree q*
shows *((λx. poly q x / poly p x) → 0) at-infinity*
 ⟨proof⟩

lemma *lead-coeff-list-def*:
lead-coeff p = (if coeffs p = [] then 0 else last (coeffs p))
 ⟨proof⟩

lemma *poly-linepath-comp*:
fixes *a::'a::{real-normed-vector,comm-semiring-0,real-algebra-1}*
shows *poly p o (linepath a b) = poly (p ∘_p [:a, b-a:]) o of-real*
 ⟨proof⟩

lemma *poly-eventually-not-zero*:
fixes *p::real poly*
assumes *p ≠ 0*
shows *eventually (λx. poly p x ≠ 0) at-infinity*
 ⟨proof⟩

2.2 More about degree

lemma *degree-div-less*:
 fixes $x y :: 'a :: \text{field}$ poly
 assumes $\text{degree } x \neq 0$ $\text{degree } y \neq 0$
 shows $\text{degree } (x \text{ div } y) < \text{degree } x$
 $\langle \text{proof} \rangle$

lemma *map-poly-degree-eq*:
 assumes $f (\text{lead-coeff } p) \neq 0$
 shows $\text{degree } (\text{map-poly } f p) = \text{degree } p$
 $\langle \text{proof} \rangle$

lemma *map-poly-degree-less*:
 assumes $f (\text{lead-coeff } p) = 0$ $\text{degree } p \neq 0$
 shows $\text{degree } (\text{map-poly } f p) < \text{degree } p$
 $\langle \text{proof} \rangle$

lemma *map-poly-degree-leq[simp]*:
 shows $\text{degree } (\text{map-poly } f p) \leq \text{degree } p$
 $\langle \text{proof} \rangle$

2.3 roots / zeros of a univariate function

definition *roots-within*::($'a \Rightarrow 'b :: \text{zero}$) $\Rightarrow 'a \text{ set} \Rightarrow 'a \text{ set}$ where
 $\text{roots-within } f s = \{x \in s. f x = 0\}$

abbreviation *roots*::($'a \Rightarrow 'b :: \text{zero}$) $\Rightarrow 'a \text{ set}$ where
 $\text{roots } f \equiv \text{roots-within } f \text{ UNIV}$

2.4 The argument principle specialised to polynomials.

lemma *argument-principle-poly*:
 assumes $p \neq 0$ and *valid*: $\text{valid-path } g$ and *loop*: $\text{pathfinish } g = \text{pathstart } g$
 and *no-proots*: $\text{path-image } g \subseteq - \text{proots } p$
 shows $\text{contour-integral } g (\lambda x. \text{deriv } (\text{poly } p) x / \text{poly } p x) = 2 * \text{of-real } \pi * i *$
 $(\sum x \in \text{proots } p. \text{winding-number } g x * \text{of-nat } (\text{order } x p))$
 $\langle \text{proof} \rangle$

end

3 Some useful lemmas about transcendental functions

theory *Missing-Transcendental* imports
 Missing-Topology
 Missing-Algebraic
begin

3.1 Misc

lemma *Im-Ln-eq-pi-half*:

$$z \neq 0 \implies (\text{Im}(\text{Ln } z) = \pi/2 \iff 0 < \text{Im}(z) \wedge \text{Re}(z) = 0)$$

$$z \neq 0 \implies (\text{Im}(\text{Ln } z) = -\pi/2 \iff \text{Im}(z) < 0 \wedge \text{Re}(z) = 0)$$

<proof>

lemma *Im-Ln-eq*:

assumes $z \neq 0$

shows $\text{Im}(\text{Ln } z) =$ (if $\text{Re } z \neq 0$ then
 if $\text{Re } z > 0$ then
 $\arctan(\text{Im } z / \text{Re } z)$
 else if $\text{Im } z \geq 0$ then
 $\arctan(\text{Im } z / \text{Re } z) + \pi$
 else
 $\arctan(\text{Im } z / \text{Re } z) - \pi$
 else
 if $\text{Im } z > 0$ then $\pi/2$ else $-\pi/2$)

<proof>

lemma *exp-Arg2pi2pi-multivalued*:

assumes $\exp(i * \text{of-real } x) = z$

shows $\exists k :: \text{int}. x = \text{Arg}2\pi z + 2 * k * \pi$

<proof>

lemma *cos-eq-neg-periodic-intro*:

assumes $x - y = 2 * (\text{of-int } k) * \pi + \pi \vee x + y = 2 * (\text{of-int } k) * \pi + \pi$

shows $\cos x = - \cos y$ *<proof>*

lemma *cos-eq-periodic-intro*:

assumes $x - y = 2 * (\text{of-int } k) * \pi \vee x + y = 2 * (\text{of-int } k) * \pi$

shows $\cos x = \cos y$ *<proof>*

lemma *sin-tan-half*: $\sin(2 * x) = 2 * \tan x / (1 + (\tan x)^2)$

<proof>

lemma *cos-tan-half*: $\cos x \neq 0 \implies \cos(2 * x) = (1 - (\tan x)^2) / (1 + (\tan x)^2)$

<proof>

lemma *tan-eq-arctan-Ex*:

shows $\tan x = y \iff (\exists k :: \text{int}. x = \arctan y + k * \pi \vee (x = \pi/2 + k * \pi \wedge y = 0))$

<proof>

lemma *arccos-unique*:

assumes $0 \leq x$

and $x \leq \pi$

and $\cos x = y$

shows $\arccos y = x$

<proof>

lemma *cos-eq-arccos-Ex*:

$\cos x = y \iff -1 \leq y \wedge y \leq 1 \wedge (\exists k::int. x = \arccos y + 2*k*\pi \vee x = -\arccos y + 2*k*\pi)$

<proof>

lemma *uniform-discrete-tan-eq*:

uniform-discrete $\{x::real. \tan x = y\}$

<proof>

lemma *get-norm-value*:

fixes $a::'a::\{\text{floor-ceiling}\}$

assumes $pp > 0$

obtains $k::int$ **and** $a1$ **where** $a = (\text{of-int } k) * pp + a1$ $a0 \leq a1$ $a1 < a0 + pp$

<proof>

lemma *filtermap-tan-at-right*:

fixes $a::real$

assumes $\cos a \neq 0$

shows $\text{filtermap } \tan (\text{at-right } a) = \text{at-right } (\tan a)$

<proof>

lemma *filtermap-tan-at-left*:

fixes $a::real$

assumes $\cos a \neq 0$

shows $\text{filtermap } \tan (\text{at-left } a) = \text{at-left } (\tan a)$

<proof>

lemma *cos-zero-iff-int2*:

fixes $x::real$

shows $\cos x = 0 \iff (\exists n::int. x = n * \pi + \pi/2)$

<proof>

lemma *filtermap-tan-at-right-inf*:

fixes $a::real$

assumes $\cos a = 0$

shows $\text{filtermap } \tan (\text{at-right } a) = \text{at-bot}$

<proof>

lemma *filtermap-tan-at-left-inf*:

fixes $a::real$

assumes $\cos a = 0$

shows $\text{filtermap } \tan (\text{at-left } a) = \text{at-top}$

<proof>

3.2 Periodic set

definition *periodic-set*:: *real set* \Rightarrow *real* \Rightarrow *bool* **where**

periodic-set $S \ \delta \longleftrightarrow (\exists B. \text{finite } B \wedge (\forall x \in S. \exists b \in B. \exists k :: \text{int}. x = b + k * \delta))$

lemma *periodic-set-multiple*:

assumes $k \neq 0$

shows *periodic-set* $S \ \delta \longleftrightarrow$ *periodic-set* $S \ (\text{of-int } k * \delta)$

<proof>

lemma *periodic-set-empty[simp]*: *periodic-set* $\{\}$ δ

<proof>

lemma *periodic-set-finite*:

assumes *finite* S

shows *periodic-set* $S \ \delta$

<proof>

lemma *periodic-set-subset[elim]*:

assumes *periodic-set* $S \ \delta \ T \subseteq S$

shows *periodic-set* $T \ \delta$

<proof>

lemma *periodic-set-union*:

assumes *periodic-set* $S \ \delta$ *periodic-set* $T \ \delta$

shows *periodic-set* $(S \cup T) \ \delta$

<proof>

lemma *periodic-imp-uniform-discrete*:

assumes *periodic-set* $S \ \delta$

shows *uniform-discrete* S

<proof>

lemma *periodic-set-tan-linear*:

assumes $a \neq 0 \ c \neq 0$

shows *periodic-set* $(\text{roots } (\lambda x. a * \tan(x/c) + b)) \ (c * \pi)$

<proof>

lemma *periodic-set-cos-linear*:

assumes $a \neq 0 \ c \neq 0$

shows *periodic-set* $(\text{roots } (\lambda x. a * \cos(x/c) + b)) \ (2 * c * \pi)$

<proof>

lemma *periodic-set-tan-poly*:

assumes $p \neq 0 \ c \neq 0$

shows *periodic-set* $(\text{roots } (\lambda x. \text{poly } p \ (\tan(x/c)))) \ (c * \pi)$

<proof>

lemma *periodic-set-sin-cos-linear*:

fixes $a \ b \ c :: \text{real}$

```

assumes  $a \neq 0 \vee b \neq 0 \vee c \neq 0$ 
shows periodic-set (roots ( $\lambda x. a * \cos x + b * \sin x + c$ )) ( $4 * \pi$ )
<proof>

end

```

4 Some useful lemmas in analysis

```

theory Missing-Analysis
imports HOL-Complex-Analysis.Complex-Analysis
begin

```

4.1 More about paths

```

lemma pathfinish-offset[simp]:
  pathfinish ( $\lambda t. g \ t - z$ ) = pathfinish  $g - z$ 
<proof>

```

```

lemma pathstart-offset[simp]:
  pathstart ( $\lambda t. g \ t - z$ ) = pathstart  $g - z$ 
<proof>

```

```

lemma pathimage-offset[simp]:
  fixes  $g :: - \Rightarrow 'b::\text{topological-group-add}$ 
  shows  $p \in \text{path-image } (\lambda t. g \ t - z) \iff p+z \in \text{path-image } g$ 
<proof>

```

```

lemma path-offset[simp]:
  fixes  $g :: - \Rightarrow 'b::\text{topological-group-add}$ 
  shows path ( $\lambda t. g \ t - z$ )  $\iff$  path  $g$ 
<proof>

```

```

lemma not-on-circlepathI:
  assumes  $c \bmod (z-z0) \neq |r|$ 
  shows  $z \notin \text{path-image } (\text{part-circlepath } z0 \ r \ st \ tt)$ 
<proof>

```

```

lemma circlepath-inj-on:
  assumes  $r > 0$ 
  shows inj-on (circlepath  $z \ r$ )  $\{0..<1\}$ 
<proof>

```

4.2 More lemmas related to winding-number

```

lemma winding-number-comp:
  assumes open  $s$  f holomorphic-on  $s$  path-image  $\gamma \subseteq s$ 
  valid-path  $\gamma \ z \notin \text{path-image } (f \circ \gamma)$ 
  shows winding-number ( $f \circ \gamma$ )  $z = 1 / (2 * \pi * i) * \text{contour-integral } \gamma (\lambda w. \text{deriv } f \ w / (f \ w - z))$ 

```

<proof>

lemma *winding-number-uminus-comp:*

assumes *valid-path* $\gamma - z \notin \text{path-image } \gamma$

shows *winding-number* $(\text{uminus} \circ \gamma) z = \text{winding-number } \gamma (-z)$

<proof>

lemma *winding-number-comp-linear:*

assumes $c \neq 0$ *valid-path* γ **and** *not-image:* $(z-b)/c \notin \text{path-image } \gamma$

shows *winding-number* $((\lambda x. c*x+b) \circ \gamma) z = \text{winding-number } \gamma ((z-b)/c)$ **(is**
 $?L = ?R)$

<proof>

end

5 Cauchy's index theorem

theory *Cauchy-Index-Theorem* **imports**

HOL-Complex-Analysis.Complex-Analysis

Sturm-Tarski.Sturm-Tarski

HOL-Computational-Algebra.Fundamental-Theorem-Algebra

Missing-Transcendental

Missing-Algebraic

Missing-Analysis

begin

This theory formalises Cauchy indices on the complex plane and relate them to winding numbers

5.1 Misc

lemma *atMostAtLeast-subset-convex:*

fixes $C :: \text{real set}$

assumes *convex* C

and $x \in C \ y \in C$

shows $\{x .. y\} \subseteq C$

<proof>

lemma *arg-elim:*

$f x \implies x = y \implies f y$

<proof>

lemma *arg-elim2:*

$f x1 \ x2 \implies x1 = y1 \implies x2 = y2 \implies f y1 \ y2$

<proof>

lemma *arg-elim3:*

$[f \ x1 \ x2 \ x3; x1 = y1; x2 = y2; x3 = y3] \implies f \ y1 \ y2 \ y3$

<proof>

lemma *IVT-strict*:

fixes $f :: 'a::\text{linear-continuum-topology} \Rightarrow 'b::\text{linorder-topology}$

assumes $(f\ a > y \wedge y > f\ b) \vee (f\ a < y \wedge y < f\ b)$ $a < b$ *continuous-on* $\{a .. b\}$
 f

shows $\exists x. a < x \wedge x < b \wedge f\ x = y$

<proof>

lemma (*in dense-linorder*) *atLeastAtMost-subseteq-greaterThanLessThan-iff*:

$\{a .. b\} \subseteq \{c <..< d\} \iff (a \leq b \implies c < a \wedge b < d)$

<proof>

lemma *Re-winding-number-half-right*:

assumes $\forall p \in \text{path-image } \gamma. \text{Re } p \geq \text{Re } z$ **and** *valid-path* γ **and** $z \notin \text{path-image}$

γ

shows $\text{Re}(\text{winding-number } \gamma\ z) = (\text{Im } (\text{Ln } (\text{pathfinish } \gamma - z)) - \text{Im } (\text{Ln } (\text{pathstart } \gamma - z))) / (2 * \pi)$

<proof>

lemma *Re-winding-number-half-upper*:

assumes $\forall p \in \text{path-image } \gamma. \text{Im } p \geq \text{Im } z$ **and** *valid-path* γ **and** $z \notin \text{path-image}$

γ

shows $\text{Re}(\text{winding-number } \gamma\ z) = (\text{Im } (\text{Ln } (i * z - i * \text{pathfinish } \gamma)) - \text{Im } (\text{Ln } (i * z - i * \text{pathstart } \gamma))) / (2 * \pi)$

<proof>

lemma *Re-winding-number-half-lower*:

assumes $\forall p \in \text{path-image } \gamma. \text{Im } p \leq \text{Im } z$ **and** *valid-path* γ **and** $z \notin \text{path-image}$

γ

shows $\text{Re}(\text{winding-number } \gamma\ z) = (\text{Im } (\text{Ln } (i * \text{pathfinish } \gamma - i * z)) - \text{Im } (\text{Ln } (i * \text{pathstart } \gamma - i * z))) / (2 * \pi)$

<proof>

lemma *Re-winding-number-half-left*:

assumes $\forall p \in \text{path-image } \gamma. \text{Re } p \leq \text{Re } z$ **and** *valid-path* γ **and** $z \notin \text{path-image}$

γ

shows $\text{Re}(\text{winding-number } \gamma\ z) = (\text{Im } (\text{Ln } (z - \text{pathfinish } \gamma)) - \text{Im } (\text{Ln } (z - \text{pathstart } \gamma))) / (2 * \pi)$

<proof>

lemma *continuous-on-open-Collect-neq*:

fixes $f\ g :: 'a::\text{topological-space} \Rightarrow 'b::\text{t2-space}$

assumes f : *continuous-on* S f **and** g : *continuous-on* S g **and** *open* S

shows *open* $\{x \in S. f\ x \neq g\ x\}$

<proof>

5.2 Sign at a filter

definition *has-sgnx*:(*real* \Rightarrow *real*) \Rightarrow *real* \Rightarrow *real filter* \Rightarrow *bool*
 (**infixr** *has'-sgnx* 55) **where**
 (*f has-sgnx c*) *F* = (*eventually* ($\lambda x. \text{sgn}(f x) = c$) *F*)

definition *sgnx-able* (**infixr** *sgnx'-able* 55) **where**
 (*f sgnx-able F*) = ($\exists c. (f \text{ has-sgnx } c) F$)

definition *sgnx* **where**
sgnx f F = (*SOME* *c. (f has-sgnx c) F*)

lemma *has-sgnx-eq-rhs*: (*f has-sgnx x*) *F* \Longrightarrow $x = y \Longrightarrow (f \text{ has-sgnx } y) F$
 \langle *proof* \rangle

named-theorems *sgnx-intros* *introduction rules for has-sgnx*
 \langle *ML* \rangle

lemma *sgnx-able-sgnx*:*f sgnx-able F* $\Longrightarrow (f \text{ has-sgnx } (\text{sgnx } f F)) F$
 \langle *proof* \rangle

lemma *has-sgnx-imp-sgnx-able[elim]*:
 (*f has-sgnx c*) *F* $\Longrightarrow f \text{ sgnx-able } F$
 \langle *proof* \rangle

lemma *has-sgnx-unique*:
assumes $F \neq \text{bot}$ (*f has-sgnx c1*) *F* (*f has-sgnx c2*) *F*
shows $c1 = c2$
 \langle *proof* \rangle

lemma *has-sgnx-imp-sgnx[elim]*:
 (*f has-sgnx c*) *F* $\Longrightarrow F \neq \text{bot} \Longrightarrow \text{sgnx } f F = c$
 \langle *proof* \rangle

lemma *has-sgnx-const[simp,sgnx-intros]*:
 ($(\lambda-. c) \text{ has-sgnx } \text{sgn } c$) *F*
 \langle *proof* \rangle

lemma *finite-sgnx-at-left-at-right*:
assumes *finite* $\{t. f t = 0 \wedge a < t \wedge t < b\}$ *continuous-on* ($\{a <..<b\} - s$) *f finite*
 s
and $x : x \in \{a <..<b\}$
shows $f \text{ sgnx-able } (\text{at-left } x) \text{ sgnx } f (\text{at-left } x) \neq 0$
 $f \text{ sgnx-able } (\text{at-right } x) \text{ sgnx } f (\text{at-right } x) \neq 0$
 \langle *proof* \rangle

lemma *sgnx-able-poly[simp]*:
 (*poly p*) *sgnx-able* (*at-right a*)
 (*poly p*) *sgnx-able* (*at-left a*)
 (*poly p*) *sgnx-able* *at-top*

(poly p) sgnx-able at-bot
<proof>

lemma *has-sgnx-identity*[intro,sgnx-intros]:
 shows $x \geq 0 \implies ((\lambda x. x) \text{ has-sgnx } 1) \text{ (at-right } x)$
 $x \leq 0 \implies ((\lambda x. x) \text{ has-sgnx } -1) \text{ (at-left } x)$
<proof>

lemma *has-sgnx-divide*[sgnx-intros]:
 assumes (f has-sgnx c1) F (g has-sgnx c2) F
 shows $((\lambda x. f x / g x) \text{ has-sgnx } c1 / c2) F$
<proof>

lemma *sgnx-able-divide*[sgnx-intros]:
 assumes f sgnx-able F g sgnx-able F
 shows $(\lambda x. f x / g x) \text{ sgnx-able } F$
<proof>

lemma *sgnx-divide*:
 assumes $F \neq \text{bot}$ f sgnx-able F g sgnx-able F
 shows $\text{sgnx } (\lambda x. f x / g x) F = \text{sgnx } f F / \text{sgnx } g F$
<proof>

lemma *has-sgnx-times*[sgnx-intros]:
 assumes (f has-sgnx c1) F (g has-sgnx c2) F
 shows $((\lambda x. f x * g x) \text{ has-sgnx } c1 * c2) F$
<proof>

lemma *sgnx-able-times*[sgnx-intros]:
 assumes f sgnx-able F g sgnx-able F
 shows $(\lambda x. f x * g x) \text{ sgnx-able } F$
<proof>

lemma *sgnx-times*:
 assumes $F \neq \text{bot}$ f sgnx-able F g sgnx-able F
 shows $\text{sgnx } (\lambda x. f x * g x) F = \text{sgnx } f F * \text{sgnx } g F$
<proof>

lemma *tendsto-nonzero-has-sgnx*:
 assumes (f \longrightarrow c) F $c \neq 0$
 shows (f has-sgnx sgn c) F
<proof>

lemma *tendsto-nonzero-sgnx*:
 assumes (f \longrightarrow c) F $F \neq \text{bot}$ $c \neq 0$
 shows $\text{sgnx } f F = \text{sgn } c$
<proof>

lemma *filterlim-divide-at-bot-at-top-iff*:

assumes $(f \longrightarrow c) F c \neq 0$

shows

$(\text{LIM } x F. f x / g x \text{ :> } \text{at-bot}) \longleftrightarrow (g \longrightarrow 0) F$
 $\wedge ((\lambda x. g x) \text{ has-sgnx } - \text{sgn } c) F$
 $(\text{LIM } x F. f x / g x \text{ :> } \text{at-top}) \longleftrightarrow (g \longrightarrow 0) F$
 $\wedge ((\lambda x. g x) \text{ has-sgnx } \text{sgn } c) F$

$\langle \text{proof} \rangle$

lemma *poly-sgnx-left-right*:

fixes $c a :: \text{real}$ **and** $p :: \text{real poly}$

assumes $p \neq 0$

shows $\text{sgnx } (\text{poly } p) (\text{at-left } a) = (\text{if even } (\text{order } a) p)$
 $\text{then } \text{sgnx } (\text{poly } p) (\text{at-right } a)$
 $\text{else } -\text{sgnx } (\text{poly } p) (\text{at-right } a)$

$\langle \text{proof} \rangle$

lemma *poly-has-sgnx-left-right*:

fixes $c a :: \text{real}$ **and** $p :: \text{real poly}$

assumes $p \neq 0$

shows $(\text{poly } p \text{ has-sgnx } c) (\text{at-left } a) \longleftrightarrow (\text{if even } (\text{order } a) p)$
 $\text{then } (\text{poly } p \text{ has-sgnx } c) (\text{at-right } a)$
 $\text{else } (\text{poly } p \text{ has-sgnx } -c) (\text{at-right } a)$

$\langle \text{proof} \rangle$

lemma *sign-r-pos-sgnx-iff*:

$\text{sign-r-pos } p a \longleftrightarrow \text{sgnx } (\text{poly } p) (\text{at-right } a) > 0$

$\langle \text{proof} \rangle$

lemma *sgnx-values*:

assumes $f \text{ sgnx-able } F F \neq \text{bot}$

shows $\text{sgnx } f F = -1 \vee \text{sgnx } f F = 0 \vee \text{sgnx } f F = 1$

$\langle \text{proof} \rangle$

lemma *has-sgnx-poly-at-top*:

$(\text{poly } p \text{ has-sgnx } \text{sgn-pos-inf } p) \text{ at-top}$

$\langle \text{proof} \rangle$

lemma *has-sgnx-poly-at-bot*:

$(\text{poly } p \text{ has-sgnx } \text{sgn-neg-inf } p) \text{ at-bot}$

$\langle \text{proof} \rangle$

lemma *sgnx-poly-at-top*:

$\text{sgnx } (\text{poly } p) \text{ at-top} = \text{sgn-pos-inf } p$

$\langle \text{proof} \rangle$

lemma *sgnx-poly-at-bot*:

$\text{sgnx } (\text{poly } p) \text{ at-bot} = \text{sgn-neg-inf } p$
 ⟨proof⟩

lemma *poly-has-sgnx-values*:

assumes $p \neq 0$

shows

$(\text{poly } p \text{ has-sgnx } 1) (\text{at-left } a) \vee (\text{poly } p \text{ has-sgnx } -1) (\text{at-left } a)$
 $(\text{poly } p \text{ has-sgnx } 1) (\text{at-right } a) \vee (\text{poly } p \text{ has-sgnx } -1) (\text{at-right } a)$
 $(\text{poly } p \text{ has-sgnx } 1) \text{ at-top} \vee (\text{poly } p \text{ has-sgnx } -1) \text{ at-top}$
 $(\text{poly } p \text{ has-sgnx } 1) \text{ at-bot} \vee (\text{poly } p \text{ has-sgnx } -1) \text{ at-bot}$

⟨proof⟩

lemma *poly-sgnx-values*:

assumes $p \neq 0$

shows $\text{sgnx } (\text{poly } p) (\text{at-left } a) = 1 \vee \text{sgnx } (\text{poly } p) (\text{at-left } a) = -1$

$\text{sgnx } (\text{poly } p) (\text{at-right } a) = 1 \vee \text{sgnx } (\text{poly } p) (\text{at-right } a) = -1$

⟨proof⟩

lemma *has-sgnx-inverse*: $(f \text{ has-sgnx } c) F \longleftrightarrow ((\text{inverse } o f) \text{ has-sgnx } (\text{inverse } c)) F$

⟨proof⟩

lemma *has-sgnx-derivative-at-left*:

assumes $g \text{-deriv}:(g \text{ has-field-derivative } c) (\text{at } x)$ and $g \ x=0$ and $c \neq 0$

shows $(g \text{ has-sgnx } - \text{sgn } c) (\text{at-left } x)$

⟨proof⟩

lemma *has-sgnx-derivative-at-right*:

assumes $g \text{-deriv}:(g \text{ has-field-derivative } c) (\text{at } x)$ and $g \ x=0$ and $c \neq 0$

shows $(g \text{ has-sgnx } \text{sgn } c) (\text{at-right } x)$

⟨proof⟩

lemma *has-sgnx-split*:

$(f \text{ has-sgnx } c) (\text{at } x) \longleftrightarrow (f \text{ has-sgnx } c) (\text{at-left } x) \wedge (f \text{ has-sgnx } c) (\text{at-right } x)$

⟨proof⟩

lemma *sgnx-at-top-IVT*:

assumes $\text{sgnx } (\text{poly } p) (\text{at-right } a) \neq \text{sgnx } (\text{poly } p) \text{ at-top}$

shows $\exists x > a. \text{poly } p \ x=0$

⟨proof⟩

lemma *sgnx-at-left-at-right-IVT*:

assumes $\text{sgnx } (\text{poly } p) (\text{at-right } a) \neq \text{sgnx } (\text{poly } p) (\text{at-left } b) \ a < b$

shows $\exists x. a < x \wedge x < b \wedge \text{poly } p \ x=0$

⟨proof⟩

lemma *sgnx-at-bot-IVT*:

assumes $\text{sgnx } (\text{poly } p) (\text{at-left } a) \neq \text{sgnx } (\text{poly } p) \text{ at-bot}$

shows $\exists x < a. \text{poly } p \ x = 0$
 <proof>

lemma *sgnx-poly-nz*:
assumes $\text{poly } p \ x \neq 0$
shows $\text{sgnx } (\text{poly } p) \ (\text{at-left } x) = \text{sgn } (\text{poly } p \ x)$
 $\text{sgnx } (\text{poly } p) \ (\text{at-right } x) = \text{sgn } (\text{poly } p \ x)$
 <proof>

5.3 Finite predicate segments over an interval

inductive *finite-Psegments*:: $(\text{real} \Rightarrow \text{bool}) \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{bool}$ for P where
emptyI: $a \geq b \Longrightarrow \text{finite-Psegments } P \ a \ b$
insertI-1: $\llbracket s \in \{a..<b\}; s = a \vee P \ s; \forall t \in \{s <..<b\}. P \ t; \text{finite-Psegments } P \ a \ s \rrbracket$
 $\Longrightarrow \text{finite-Psegments } P \ a \ b$
insertI-2: $\llbracket s \in \{a..<b\}; s = a \vee P \ s; (\forall t \in \{s <..<b\}. \neg P \ t); \text{finite-Psegments } P \ a \ s \rrbracket$
 $\Longrightarrow \text{finite-Psegments } P \ a \ b$

lemma *finite-Psegments-pos-linear*:
assumes $\text{finite-Psegments } P \ (b * lb + c) \ (b * ub + c)$ and $b > 0$
shows $\text{finite-Psegments } (P \ o \ (\lambda t. b * t + c)) \ lb \ ub$
 <proof>

lemma *finite-Psegments-congE*:
assumes $\text{finite-Psegments } Q \ lb \ ub$
 $\bigwedge t. \llbracket lb < t; t < ub \rrbracket \Longrightarrow Q \ t \longleftrightarrow P \ t$
shows $\text{finite-Psegments } P \ lb \ ub$ <proof>

lemma *finite-Psegments-constI*:
assumes $\bigwedge t. \llbracket a < t; t < b \rrbracket \Longrightarrow P \ t = c$
shows $\text{finite-Psegments } P \ a \ b$
 <proof>

context
begin

private lemma *finite-Psegments-less-eq1*:
assumes $\text{finite-Psegments } P \ a \ c \ b \leq c$
shows $\text{finite-Psegments } P \ a \ b$ <proof> **lemma** *finite-Psegments-less-eq2*:
assumes $\text{finite-Psegments } P \ a \ c \ a \leq b$
shows $\text{finite-Psegments } P \ b \ c$ <proof>

lemma *finite-Psegments-included*:
assumes $\text{finite-Psegments } P \ a \ d \ a \leq b \ c \leq d$
shows $\text{finite-Psegments } P \ b \ c$
 <proof>
end

lemma *finite-Psegments-combine*:

assumes *finite-Psegments* P a b *finite-Psegments* P b c $b \in \{a..c\}$ *closed* ($\{x. P\ x\} \cap \{a..c\}$)
shows *finite-Psegments* P a c *<proof>*

5.4 Finite segment intersection of a path with the imaginary axis

definition *finite-ReZ-segments::*(*real* \Rightarrow *complex*) \Rightarrow *complex* \Rightarrow *bool* **where**

finite-ReZ-segments g $z = \text{finite-Psegments } (\lambda t. \text{Re } (g\ t - z) = 0)\ 0\ 1$

lemma *finite-ReZ-segments-joinpaths*:

assumes $g1$:*finite-ReZ-segments* $g1$ z **and** $g2$: *finite-ReZ-segments* $g2$ z **and**
 $\text{path } g1\ \text{path } g2\ \text{pathfinish } g1 = \text{pathstart } g2$
shows *finite-ReZ-segments* $(g1+++g2)$ z
<proof>

lemma *finite-ReZ-segments-congE*:

assumes *finite-ReZ-segments* $p1$ $z1$
 $\bigwedge t. \llbracket 0 < t; t < 1 \rrbracket \implies \text{Re}(p1\ t - z1) = \text{Re}(p2\ t - z2)$
shows *finite-ReZ-segments* $p2$ $z2$
<proof>

lemma *finite-ReZ-segments-constI*:

assumes $\forall t. 0 < t \wedge t < 1 \longrightarrow g\ t = c$
shows *finite-ReZ-segments* g z
<proof>

lemma *finite-ReZ-segment-cases* [*consumes 1, case-names subEq subNEq, cases pred:finite-ReZ-segments*]:

assumes *finite-ReZ-segments* g z
and *subEq*:($\bigwedge s. \llbracket s \in \{0..<1\}; s=0 \vee \text{Re } (g\ s) = \text{Re } z; \forall t \in \{s < .. < 1\}. \text{Re } (g\ t) = \text{Re } z; \text{finite-ReZ-segments } (\text{subpath } 0\ s\ g)\ z \rrbracket \implies P$)
and *subNEq*:($\bigwedge s. \llbracket s \in \{0..<1\}; s=0 \vee \text{Re } (g\ s) = \text{Re } z; \forall t \in \{s < .. < 1\}. \text{Re } (g\ t) \neq \text{Re } z; \text{finite-ReZ-segments } (\text{subpath } 0\ s\ g)\ z \rrbracket \implies P$)
shows P
<proof>

lemma *finite-ReZ-segments-induct* [*case-names sub0 subEq subNEq, induct pred:finite-ReZ-segments*]:

assumes *finite-ReZ-segments* g z
assumes *sub0*: $\bigwedge g\ z. (P\ (\text{subpath } 0\ 0\ g)\ z)$
and *subEq*:($\bigwedge s\ g\ z. \llbracket s \in \{0..<1\}; s=0 \vee \text{Re } (g\ s) = \text{Re } z; \forall t \in \{s < .. < 1\}. \text{Re } (g\ t) = \text{Re } z; \text{finite-ReZ-segments } (\text{subpath } 0\ s\ g)\ z; P\ (\text{subpath } 0\ s\ g)\ z \rrbracket \implies P\ g\ z$)
and *subNEq*:($\bigwedge s\ g\ z. \llbracket s \in \{0..<1\}; s=0 \vee \text{Re } (g\ s) = \text{Re } z; \forall t \in \{s < .. < 1\}. \text{Re } (g\ t) \neq \text{Re } z; \text{finite-ReZ-segments } (\text{subpath } 0\ s\ g)\ z; P\ (\text{subpath } 0\ s\ g)\ z \rrbracket \implies P\ g\ z$)

shows $P g z$
<proof>

lemma *finite-ReZ-segments-shiftpah:*

assumes *finite-ReZ-segments* $g z s \in \{0..1\}$ **path** g **and** *loop:pathfinish* $g = \text{path-start } g$

shows *finite-ReZ-segments* (*shiftpath* $s g$) z
<proof>

lemma *finite-imp-finite-ReZ-segments:*

assumes *finite* $\{t. \text{Re } (g t - z) = 0 \wedge 0 \leq t \wedge t \leq 1\}$

shows *finite-ReZ-segments* $g z$
<proof>

lemma *finite-ReZ-segments-poly-linepath:*

shows *finite-ReZ-segments* (*poly* $p o \text{linepath } a b$) z
<proof>

lemma *part-circlepath-half-finite-inter:*

assumes $st \neq tt \ r \neq 0 \ c \neq 0$

shows *finite* $\{t. \text{part-circlepath } z0 r st tt t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$ (**is** *finite* $?T$)
<proof>

lemma *linepath-half-finite-inter:*

assumes $a \cdot c \neq d \vee b \cdot c \neq d$

shows *finite* $\{t. \text{linepath } a b t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$ (**is** *finite* $?S$)
<proof>

lemma *finite-half-joinpaths-inter:*

assumes *finite* $\{t. l1 t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$ *finite* $\{t. l2 t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$

shows *finite* $\{t. (l1+++l2) t \cdot c = d \wedge 0 \leq t \wedge t \leq 1\}$
<proof>

lemma *finite-ReZ-segments-linepath:*

finite-ReZ-segments (*linepath* $a b$) z
<proof>

lemma *finite-ReZ-segments-part-circlepath:*

finite-ReZ-segments (*part-circlepath* $z0 r st tt$) z
<proof>

lemma *finite-ReZ-segments-poly-of-real:*

shows *finite-ReZ-segments* (*poly* $p o \text{of-real}$) z

<proof>

lemma *finite-ReZ-segments-subpath:*

assumes *finite-ReZ-segments* $g z$

$0 \leq u \ u \leq v \ v \leq 1$

shows *finite-ReZ-segments* (*subpath u v g*) *z*
 ⟨*proof*⟩

5.5 jump and jumpF

definition *jump*::(*real* ⇒ *real*) ⇒ *real* ⇒ *int* **where**

jump f a = (if
 (*LIM x (at-left a). f x* :> *at-bot*) ∧ (*LIM x (at-right a). f x* :> *at-top*)
 then 1 else if
 (*LIM x (at-left a). f x* :> *at-top*) ∧ (*LIM x (at-right a). f x* :> *at-bot*)
 then -1 else 0)

definition *jumpF*::(*real* ⇒ *real*) ⇒ *real filter* ⇒ *real* **where**

jumpF f F ≡ (if *filterlim f at-top F* then 1/2 else
 if *filterlim f at-bot F* then -1/2 else (0::*real*))

lemma *jumpF-const[simp]*:

assumes *F ≠ bot*
shows *jumpF (λ-. c) F* = 0

⟨*proof*⟩

lemma *jumpF-not-infinity*:

assumes *continuous F g F ≠ bot*
shows *jumpF g F* = 0

⟨*proof*⟩

lemma *jumpF-linear-comp*:

assumes *c ≠ 0*

shows

*jumpF (f o (λx. c*x+b)) (at-left x)* =
 (if *c > 0* then *jumpF f (at-left (c*x+b))* else *jumpF f (at-right (c*x+b))*)
 (is ?*case1*)
*jumpF (f o (λx. c*x+b)) (at-right x)* =
 (if *c > 0* then *jumpF f (at-right (c*x+b))* else *jumpF f (at-left (c*x+b))*)
 (is ?*case2*)

⟨*proof*⟩

lemma *jump-const[simp]:jump (λ-. c) a* = 0

⟨*proof*⟩

lemma *jump-not-infinity*:

isCont f a ⇒ *jump f a* = 0

⟨*proof*⟩

lemma *jump-jump-poly-aux*:

assumes *p ≠ 0 coprime p q*

shows *jump (λx. poly q x / poly p x) a* = *jump-poly q p a*

⟨*proof*⟩

lemma *jump-jumpF*:

assumes *cont:isCont* (*inverse o f*) *a* **and**

sgn_l:(*f has-sgn_l* *l*) (*at-left a*) **and** *sgn_r*:(*f has-sgn_r* *r*) (*at-right a*) **and**
l≠0 *r*≠0

shows *jump f a* = *jumpF f (at-right a)* - *jumpF f (at-left a)*

⟨*proof*⟩

lemma *jump-linear-comp*:

assumes *c*≠0

shows *jump (f o (λx. c*x+b)) x* = (*if c*>0 *then jump f (c*x+b)* *else -jump f*
*(c*x+b)*)

⟨*proof*⟩

lemma *jump-divide-derivative*:

assumes *isCont f x g x* = 0 *f x*≠0

and *g-deriv*:(*g has-field-derivative c*) (*at x*) **and** *c*≠0

shows *jump (λt. f t/g t) x* = (*if sgn c* = *sgn (f x)* *then 1* *else -1*)

⟨*proof*⟩

lemma *jump-jump-poly*: *jump (λx. poly q x / poly p x) a* = *jump-poly q p a*

⟨*proof*⟩

lemma *jump-Im-divide-Re-0*:

assumes *path g Re (g x)*≠0 *0*<*x* <*1*

shows *jump (λt. Im (g t) / Re (g t)) x* = 0

⟨*proof*⟩

lemma *jumpF-im-divide-Re-0*:

assumes *path g Re (g x)*≠0

shows $\llbracket 0 \leq x; x < 1 \rrbracket \implies \text{jumpF } (\lambda t. \text{Im } (g t) / \text{Re } (g t)) \text{ (at-right } x) = 0$

$\llbracket 0 < x; x \leq 1 \rrbracket \implies \text{jumpF } (\lambda t. \text{Im } (g t) / \text{Re } (g t)) \text{ (at-left } x) = 0$

⟨*proof*⟩

lemma *jump-cong*:

assumes *x*=*y* **and** *eventually* (*λx. f x=g x*) (*at x*)

shows *jump f x* = *jump g y*

⟨*proof*⟩

lemma *jumpF-cong*:

assumes *F*=*G* **and** *eventually* (*λx. f x=g x*) *F*

shows *jumpF f F* = *jumpF g G*

⟨*proof*⟩

lemma *jump-at-left-at-right-eq*:

assumes *isCont f x* **and** *f x* ≠ 0 **and** *sgn_l-eq*:*sgn_l g (at-left x)* = *sgn_r g (at-right*
x)

shows *jump (λt. f t/g t) x* = 0

⟨*proof*⟩

lemma *jumpF-pos-has-sgnx*:

assumes *jumpF f F > 0*

shows (*f has-sgnx 1*) *F*

<proof>

lemma *jumpF-neg-has-sgnx*:

assumes *jumpF f F < 0*

shows (*f has-sgnx -1*) *F*

<proof>

lemma *jumpF-IVT*:

fixes *f::real ⇒ real* **and** *a b::real*

defines *right*≡($\lambda(R::real \Rightarrow real \Rightarrow bool). R (jumpF f (at-right a)) 0$
 $\vee (continuous (at-right a) f \wedge R (f a) 0)$)

and

left≡($\lambda(R::real \Rightarrow real \Rightarrow bool). R (jumpF f (at-left b)) 0$
 $\vee (continuous (at-left b) f \wedge R (f b) 0)$)

assumes *a < b* **and** *cont:continuous-on {a<..**<b**}* *f* **and**

right-left:right greater \wedge *left less* \vee *right less* \wedge *left greater*

shows $\exists x. a < x \wedge x < b \wedge f x = 0$

<proof>

lemma *jumpF-eventually-const*:

assumes *eventually* ($\lambda x. f x = c$) *F F ≠ bot*

shows *jumpF f F = 0*

<proof>

lemma *jumpF-tan-comp*:

jumpF (f o tan) (at-right x) = (if cos x = 0
then jumpF f at-bot else jumpF f (at-right (tan x)))

jumpF (f o tan) (at-left x) = (if cos x = 0
then jumpF f at-top else jumpF f (at-left (tan x)))

<proof>

5.6 Finite jumpFs over an interval

definition *finite-jumpFs::(real ⇒ real) ⇒ real ⇒ real ⇒ bool* **where**

finite-jumpFs f a b = finite {x. (jumpF f (at-left x) ≠ 0 \vee *jumpF f (at-right x)*
≠ 0) \wedge a ≤ x \wedge x ≤ b}

lemma *finite-jumpFs-linear-pos*:

assumes *c > 0*

shows *finite-jumpFs (f o (λx. c * x + b)) lb ub* \longleftrightarrow *finite-jumpFs f (c * lb + b)*
*(c * ub + b)*

<proof>

lemma *finite-jumpFs-consts*:

finite-jumpFs ($\lambda \cdot c$) lb ub
<proof>

lemma *finite-jumpFs-combine*:
 assumes *finite-jumpFs* f a b *finite-jumpFs* f b c
 shows *finite-jumpFs* f a c
<proof>

lemma *finite-jumpFs-subE*:
 assumes *finite-jumpFs* f a b $a \leq a'$ $b' \leq b$
 shows *finite-jumpFs* f a' b'
<proof>

lemma *finite-Psegments-Re-imp-jumpFs*:
 assumes *finite-Psegments* ($\lambda t. \text{Re}(g t - z) = 0$) a b *continuous-on* {a..b} g
 shows *finite-jumpFs* ($\lambda t. \text{Im}(g t - z)/\text{Re}(g t - z)$) a b
<proof>

lemma *finite-ReZ-segments-imp-jumpFs*:
 assumes *finite-ReZ-segments* g z *path* g
 shows *finite-jumpFs* ($\lambda t. \text{Im}(g t - z)/\text{Re}(g t - z)$) 0 1
<proof>

5.7 *jumpF* at path ends

definition *jumpF-pathstart*::(*real* \Rightarrow *complex*) \Rightarrow *complex* \Rightarrow *real* **where**
 jumpF-pathstart g z = *jumpF* ($\lambda t. \text{Im}(g t - z)/\text{Re}(g t - z)$) (*at-right* 0)

definition *jumpF-pathfinish*::(*real* \Rightarrow *complex*) \Rightarrow *complex* \Rightarrow *real* **where**
 jumpF-pathfinish g z = *jumpF* ($\lambda t. \text{Im}(g t - z)/\text{Re}(g t - z)$) (*at-left* 1)

lemma *jumpF-pathstart-eq-0*:
 assumes *path* g $\text{Re}(\text{pathstart } g) \neq \text{Re } z$
 shows *jumpF-pathstart* g z = 0
<proof>

lemma *jumpF-pathfinish-eq-0*:
 assumes *path* g $\text{Re}(\text{pathfinish } g) \neq \text{Re } z$
 shows *jumpF-pathfinish* g z = 0
<proof>

lemma
 shows *jumpF-pathfinish-reversepath*: *jumpF-pathfinish* (*reversepath* g) z = *jumpF-pathstart*
 g z
 and *jumpF-pathstart-reversepath*: *jumpF-pathstart* (*reversepath* g) z = *jumpF-pathfinish*
 g z
<proof>

lemma *jumpF-pathstart-joinpaths[simp]*:

jumpF-pathstart (g1+++g2) z = *jumpF-pathstart* g1 z
 ⟨proof⟩

lemma *jumpF-pathfinish-joinpaths[simp]*:
jumpF-pathfinish (g1+++g2) z = *jumpF-pathfinish* g2 z
 ⟨proof⟩

5.8 Cauchy index

— Deprecated, use "cindexE" if possible

definition *cindex::real ⇒ real ⇒ (real ⇒ real) ⇒ int* **where**
cindex a b f = (∑ x∈{x. *jump* f x ≠ 0 ∧ a < x ∧ x < b}. *jump* f x)

definition *cindexE::real ⇒ real ⇒ (real ⇒ real) ⇒ real* **where**
cindexE a b f = (∑ x∈{x. *jumpF* f (at-right x) ≠ 0 ∧ a ≤ x ∧ x < b}. *jumpF* f (at-right x))
 − (∑ x∈{x. *jumpF* f (at-left x) ≠ 0 ∧ a < x ∧ x ≤ b}. *jumpF* f (at-left x))

definition *cindexE-ubd::(real ⇒ real) ⇒ real* **where**
cindexE-ubd f = (∑ x∈{x. *jumpF* f (at-right x) ≠ 0 } . *jumpF* f (at-right x))
 − (∑ x∈{x. *jumpF* f (at-left x) ≠ 0 } . *jumpF* f (at-left x))

lemma *cindexE-empty*:
cindexE a a f = 0
 ⟨proof⟩

lemma *cindex-const*: *cindex* a b (λ-. c) = 0
 ⟨proof⟩

lemma *cindex-eq-cindex-poly*: *cindex* a b (λx. poly q x / poly p x) = *cindex-poly* a b q p
 ⟨proof⟩

lemma *cindex-combine*:
assumes *finite:finite* {x. *jump* f x ≠ 0 ∧ a < x ∧ x < c} **and** a < b b < c
shows *cindex* a c f = *cindex* a b f + *jump* f b + *cindex* b c f
 ⟨proof⟩

lemma *cindexE-combine*:
assumes *finite:finite-jumpFs* f a c **and** a ≤ b b ≤ c
shows *cindexE* a c f = *cindexE* a b f + *cindexE* b c f
 ⟨proof⟩

lemma *cindex-linear-comp*:
assumes c ≠ 0
shows *cindex* lb ub (f o (λx. c*x+b)) = (if c > 0
 then *cindex* (c*lb+b) (c*ub+b) f

else - $cindex (c*ub+b) (c*lb+b) f$
 <proof>

lemma *cindexE-linear-comp*:

assumes $c \neq 0$
shows $cindexE lb ub (f o (\lambda x. c*x+b)) = (if c > 0$
 then $cindexE (c*lb+b) (c*ub+b) f$
 else - $cindexE (c*ub+b) (c*lb+b) f$)
 <proof>

lemma *cindexE-cong*:

assumes *finite* s **and** $fg\text{-}eq: \bigwedge x. \llbracket a < x; x < b; x \notin s \rrbracket \implies f x = g x$
shows $cindexE a b f = cindexE a b g$
 <proof>

lemma *cindexE-constI*:

assumes $\bigwedge t. \llbracket a < t; t < b \rrbracket \implies f t = c$
shows $cindexE a b f = 0$
 <proof>

lemma *cindex-eq-cindexE-divide*:

fixes $f g: real \Rightarrow real$
defines $h \equiv (\lambda x. f x / g x)$
assumes $a < b$ **and**
 $finite\text{-}fg: finite \{x. (f x = 0 \vee g x = 0) \wedge a \leq x \wedge x \leq b\}$ **and**
 $g\text{-}imp\text{-}f: \forall x \in \{a..b\}. g x = 0 \longrightarrow f x \neq 0$ **and**
 $f\text{-}cont: continuous\text{-}on \{a..b\} f$ **and**
 $g\text{-}cont: continuous\text{-}on \{a..b\} g$
shows $cindexE a b h = jumpF h (at\text{-}right a) + cindex a b h - jumpF h (at\text{-}left b)$
 <proof>

5.9 Cauchy index along a path

— Deprecated, use "cindex_pathE" if possible

definition *cindex-path*:: $(real \Rightarrow complex) \Rightarrow complex \Rightarrow int$ **where**
 $cindex\text{-}path g z = cindex 0 1 (\lambda t. Im (g t - z) / Re (g t - z))$

definition *cindex-pathE*:: $(real \Rightarrow complex) \Rightarrow complex \Rightarrow real$ **where**
 $cindex\text{-}pathE g z = cindexE 0 1 (\lambda t. Im (g t - z) / Re (g t - z))$

lemma *cindex-pathE-point*: $cindex\text{-}pathE (linepath a a) b = 0$
 <proof>

lemma *cindex-path-reversepath*:

$cindex\text{-}path (reversepath g) z = - cindex\text{-}path g z$
 <proof>

lemma *cindex-pathE-reversepath*: $cindex\text{-}pathE (reversepath g) z = - cindex\text{-}pathE$

$g z$
 ⟨proof⟩

lemma *cindex-pathE-reversepath'*: $cindex-pathE g z = -cindex-pathE (reversepath g) z$
 ⟨proof⟩

lemma *cindex-pathE-joinpaths*:
assumes $g1:finite-ReZ-segments g1 z$ **and** $g2:finite-ReZ-segments g2 z$ **and**
 $path g1 path g2 pathfinish g1 = pathstart g2$
shows $cindex-pathE (g1+++g2) z = cindex-pathE g1 z + cindex-pathE g2 z$
 ⟨proof⟩

lemma *cindex-pathE-constI*:
assumes $\bigwedge t. \llbracket 0 < t; t < 1 \rrbracket \implies g t = c$
shows $cindex-pathE g z = 0$
 ⟨proof⟩

lemma *cindex-pathE-subpath-combine*:
assumes $g:finite-ReZ-segments g z$ **and** $path g$ **and**
 $0 \leq a \leq b \leq c \leq 1$
shows $cindex-pathE (subpath a b g) z + cindex-pathE (subpath b c g) z$
 $= cindex-pathE (subpath a c g) z$
 ⟨proof⟩

lemma *cindex-pathE-shiftpath*:
assumes $finite-ReZ-segments g z$ $s \in \{0..1\}$ $path g$ **and** $loop:pathfinish g = pathstart g$
shows $cindex-pathE (shiftpath s g) z = cindex-pathE g z$
 ⟨proof⟩

5.10 Cauchy's Index Theorem

theorem *winding-number-cindex-pathE-aux*:
fixes $g::real \Rightarrow complex$
assumes $finite-ReZ-segments g z$ **and** $valid-path g z \notin path-image g$ **and**
 $Re-ends:Re (g 1) = Re z$ $Re (g 0) = Re z$
shows $2 * Re(winding-number g z) = - cindex-pathE g z$
 ⟨proof⟩

theorem *winding-number-cindex-pathE*:
fixes $g::real \Rightarrow complex$
assumes $finite-ReZ-segments g z$ **and** $valid-path g z \notin path-image g$ **and**
 $loop: pathfinish g = pathstart g$
shows $winding-number g z = - cindex-pathE g z / 2$
 ⟨proof⟩

REMARK: The usual statement of Cauchy's Index theorem (i.e. Analytic Theory of Polynomials (2002): Theorem 11.1.3) is about the equality between the number of polynomial roots and the Cauchy index, which

is the joint application of $\llbracket \text{finite-ReZ-segments } ?g \ ?z; \text{ valid-path } ?g; ?z \notin \text{path-image } ?g; \text{ pathfinish } ?g = \text{pathstart } ?g \rrbracket \implies \text{winding-number } ?g \ ?z = \text{complex-of-real } (- \text{cindex-pathE } ?g \ ?z / 2)$ and $\llbracket \text{open } ?s; \text{ connected } ?s; ?f \text{ holomorphic-on } ?s - ?poles; ?h \text{ holomorphic-on } ?s; \text{ valid-path } ?g; \text{ pathfinish } ?g = \text{pathstart } ?g; \text{ path-image } ?g \subseteq ?s - \{w. ?f \ w = 0 \vee w \in ?poles\}; \forall z. z \notin ?s \longrightarrow \text{winding-number } ?g \ z = 0; \text{ finite } \{w. ?f \ w = 0 \vee w \in ?poles\}; \forall p \in ?poles. \text{ is-pole } ?f \ p \rrbracket \implies \text{contour-integral } ?g \ (\lambda x. \text{deriv } ?f \ x * ?h \ x / ?f \ x) = \text{complex-of-real } (2 * \text{pi}) * \text{i} * (\sum p \in \{w. ?f \ w = 0 \vee w \in ?poles\}. \text{winding-number } ?g \ p * ?h \ p * \text{of-int } (zorder \ ?f \ p)).$

end

6 Evaluate winding numbers by calculating Cauchy indices

theory *Winding-Number-Eval* imports

Cauchy-Index-Theorem

HOL-Eisbach.Eisbach-Tools

begin

6.1 Misc

lemma *not-on-closed-segmentI*:

fixes $z::'a::\text{euclidean-space}$

assumes $\text{norm } (z - a) *_R (b - z) \neq \text{norm } (b - z) *_R (z - a)$

shows $z \notin \text{closed-segment } a \ b$

<proof>

lemma *not-on-closed-segmentI-complex*:

fixes $z::\text{complex}$

assumes $(\text{Re } b - \text{Re } z) * (\text{Im } z - \text{Im } a) \neq (\text{Im } b - \text{Im } z) * (\text{Re } z - \text{Re } a)$

shows $z \notin \text{closed-segment } a \ b$

<proof>

6.2 finite intersection with the two axes

definition *finite-axes-cross*:: $(\text{real} \Rightarrow \text{complex}) \Rightarrow \text{complex} \Rightarrow \text{bool}$ **where**

$\text{finite-axes-cross } g \ z = \text{finite } \{t. (\text{Re } (g \ t - z) = 0 \vee \text{Im } (g \ t - z) = 0) \wedge 0 \leq t \wedge t \leq 1\}$

lemma *finite-cross-intros*:

$\llbracket \text{Re } a \neq \text{Re } z \vee \text{Re } b \neq \text{Re } z; \text{Im } a \neq \text{Im } z \vee \text{Im } b \neq \text{Im } z \rrbracket \implies \text{finite-axes-cross } (\text{linepath } a \ b) \ z$

$\llbracket st \neq tt; r \neq 0 \rrbracket \implies \text{finite-axes-cross } (\text{part-circlepath } z0 \ r \ st \ tt) \ z$

$\llbracket \text{finite-axes-cross } g1 \ z; \text{finite-axes-cross } g2 \ z \rrbracket \implies \text{finite-axes-cross } (g1 +++ g2) \ z$

<proof>

lemma *cindex-path-joinpaths*:

assumes *finite-axes-cross* $g1\ z\ finite-axes-cross\ g2\ z$
and *path* $g1\ path\ g2\ pathfinish\ g1 = pathstart\ g2\ pathfinish\ g1 \neq z$
shows *cindex-path* $(g1+++g2)\ z = cindex-path\ g1\ z + jumpF-pathstart\ g2\ z$
 $- jumpF-pathfinish\ g1\ z + cindex-path\ g2\ z$
 ⟨*proof*⟩

6.3 More lemmas related *cindex-pathE* / *jumpF-pathstart* / *jumpF-pathfinish*

lemma *cindex-pathE-linepath*:

assumes $z \notin closed-segment\ a\ b$
shows *cindex-pathE* $(linepath\ a\ b)\ z =$
 $let\ c1 = Re\ a - Re\ z;$
 $c2 = Re\ b - Re\ z;$
 $c3 = Im\ a * Re\ b + Re\ z * Im\ b + Im\ z * Re\ a - Im\ z * Re\ b - Im\ b * Re\ a - Re\ z * Im\ a;$
 $d1 = Im\ a - Im\ z;$
 $d2 = Im\ b - Im\ z$
 in if $(c1 > 0 \wedge c2 < 0) \vee (c1 < 0 \wedge c2 > 0)$ then
 (if $c3 > 0$ then 1 else -1)
 else
 (if $(c1 = 0 \iff c2 \neq 0) \wedge (c1 = 0 \implies d1 \neq 0) \wedge (c2 = 0 \implies d2 \neq 0)$ then
 if $(c1 = 0 \wedge (c2 > 0 \iff d1 > 0)) \vee (c2 = 0 \wedge (c1 > 0 \iff d2 < 0))$ then
 $1/2$ else $-1/2$
 else 0)
 ⟨*proof*⟩

lemma *cindex-path-linepath*:

assumes $z \notin path-image\ (linepath\ a\ b)$
shows *cindex-path* $(linepath\ a\ b)\ z =$
 $let\ c1 = Re(a) - Re(z); c2 = Re(b) - Re(z);$
 $c3 = Im(a) * Re(b) + Re(z) * Im(b) + Im(z) * Re(a) - Im(z) * Re(b) - Im(b) * Re(a)$
 $- Re(z) * Im(a)$
 in if $(c1 > 0 \wedge c2 < 0) \vee (c1 < 0 \wedge c2 > 0)$ then (if $c3 > 0$ then 1 else -1) else
 0)
 ⟨*proof*⟩

lemma *cindex-pathE-part-circlepath*:

assumes $cmod\ (z - z0) \neq r$ **and** $r > 0\ 0 \leq st\ st < tt\ tt \leq 2 * \pi$
shows *cindex-pathE* $(part-circlepath\ z\ r\ st\ tt)\ z0 =$
 if $|Re\ z - Re\ z0| < r$ then
 (let
 $\vartheta = arccos\ ((Re\ z0 - Re\ z)/r);$
 $\beta = 2 * \pi - \vartheta$
 in
 $jumpF-pathstart\ (part-circlepath\ z\ r\ st\ tt)\ z0$
 $+$
 (if $st < \vartheta \wedge \vartheta < tt$ then if $r * \sin\ \vartheta + Im\ z > Im\ z0$ then -1 else 1 else 0)
 $+$
 (if $st < \beta \wedge \beta < tt$ then if $r * \sin\ \beta + Im\ z > Im\ z0$ then 1 else -1 else 0)

```

)
jumpF-pathfinish (part-circlepath z r st tt) z0
)
else
  if |Re z - Re z0| = r then
    jumpF-pathstart (part-circlepath z r st tt) z0
    - jumpF-pathfinish (part-circlepath z r st tt) z0
  else 0
)
⟨proof⟩

```

lemma *jumpF-pathstart-part-circlepath*:
assumes $st < tt$ $r > 0$ $cmod (z - z0) \neq r$
shows $jumpF-pathstart (part-circlepath z r st tt) z0 =$ (
 if $r * \cos st + Re z - Re z0 = 0$ then
 (let
 $\Delta = r * \sin st + Im z - Im z0$
 in
 if $(\sin st > 0 \vee \cos st = 1) \wedge \Delta < 0$
 $\vee (\sin st < 0 \vee \cos st = -1) \wedge \Delta > 0$ then
 $1/2$
 else
 $- 1/2$)
 else 0)
 ⟨proof⟩

lemma *jumpF-pathfinish-part-circlepath*:
assumes $st < tt$ $r > 0$ $cmod (z - z0) \neq r$
shows $jumpF-pathfinish (part-circlepath z r st tt) z0 =$ (
 if $r * \cos tt + Re z - Re z0 = 0$ then
 (let
 $\Delta = r * \sin tt + Im z - Im z0$
 in
 if $(\sin tt > 0 \vee \cos tt = -1) \wedge \Delta < 0$
 $\vee (\sin tt < 0 \vee \cos tt = 1) \wedge \Delta > 0$ then
 $- 1/2$
 else
 $1/2$)
 else 0)
 ⟨proof⟩

lemma
fixes $z0 z :: complex$ **and** $r :: real$
defines $upper \equiv cindex-pathE (part-circlepath z r 0 pi) z0$
and $lower \equiv cindex-pathE (part-circlepath z r pi (2*pi)) z0$
shows *cindex-pathE-circlepath-upper*:
 $\llbracket cmod (z0 - z) < r \rrbracket \implies upper = -1$
 $\llbracket Im (z0 - z) > r; |Re (z0 - z)| < r \rrbracket \implies upper = 1$
 $\llbracket Im (z0 - z) < -r; |Re (z0 - z)| < r \rrbracket \implies upper = -1$

$\llbracket \text{Re } (z0 - z) \rrbracket > r; r > 0 \rrbracket \implies \text{upper} = 0$
and *cindex-pathE-circlepath-lower*:
 $\llbracket \text{cmod } (z0 - z) < r \rrbracket \implies \text{lower} = -1$
 $\llbracket \text{Im } (z0 - z) > r; |\text{Re } (z0 - z)| < r \rrbracket \implies \text{lower} = -1$
 $\llbracket \text{Im } (z0 - z) < -r; |\text{Re } (z0 - z)| < r \rrbracket \implies \text{lower} = 1$
 $\llbracket \text{Re } (z0 - z) \rrbracket > r; r > 0 \rrbracket \implies \text{lower} = 0$
 <proof>

lemma *jumpF-pathstart-linepath*:
jumpF-pathstart (linepath a b) z =
 (if Re a = Re z \wedge Im a \neq Im z \wedge Re b \neq Re a then
 if (Im a > Im z \wedge Re b > Re a) \vee (Im a < Im z \wedge Re b < Re a) then 1/2
 else -1/2
 else 0)
 <proof>

lemma *jumpF-pathfinish-linepath*:
jumpF-pathfinish (linepath a b) z =
 (if Re b = Re z \wedge Im b \neq Im z \wedge Re a \neq Re a then
 if (Im b > Im z \wedge Re a > Re b) \vee (Im b < Im z \wedge Re a < Re b) then 1/2
 else -1/2
 else 0)
 <proof>

6.4 Setting up the method for evaluating winding numbers

lemma *pathfinish-pathstart-partcirclepath-simps*:
pathstart (part-circlepath z0 r (3*pi/2) tt) = z0 - Complex 0 r
pathstart (part-circlepath z0 r (2*pi) tt) = z0 + r
pathfinish (part-circlepath z0 r st (3*pi/2)) = z0 - Complex 0 r
pathfinish (part-circlepath z0 r st (2*pi)) = z0 + r
pathstart (part-circlepath z0 r 0 tt) = z0 + r
pathstart (part-circlepath z0 r (pi/2) tt) = z0 + Complex 0 r
pathstart (part-circlepath z0 r (pi) tt) = z0 - r
pathfinish (part-circlepath z0 r st 0) = z0 + r
pathfinish (part-circlepath z0 r st (pi/2)) = z0 + Complex 0 r
pathfinish (part-circlepath z0 r st (pi)) = z0 - r
 <proof>

lemma *winding-eq-intro*:
finite-ReZ-segments g z \implies
valid-path g \implies
 z \notin path-image g \implies
pathfinish g = *pathstart* g \implies
 - of-real (cindex-pathE g z) = 2*n \implies
winding-number g z = (n::complex)
 <proof>

named-theorems *winding-intros* and *winding-simps*

```

lemmas [winding-intros] =
  finite-ReZ-segments-joinpaths
  valid-path-join
  path-join-imp
  not-in-path-image-join

```

```

lemmas [winding-simps] =
  finite-ReZ-segments-linepath
  finite-ReZ-segments-part-circlepath
  jumpF-pathfinish-joinpaths
  jumpF-pathstart-joinpaths
  pathfinish-linepath
  pathstart-linepath
  pathfinish-join
  pathstart-join
  valid-path-linepath
  valid-path-part-circlepath
  path-part-circlepath
  Re-complex-of-real
  Im-complex-of-real
  of-real-linepath
  pathfinish-pathstart-partcirclepath-simps

```

```

method rep-subst =
  (subst cindex-pathE-joinpaths; rep-subst)?

```

The method "eval_winding" $1::'a$ will try to simplify of the form *winding-number* $g z = n$ where n is an integer and g is a closed path comprised of *linepath*, *part-circlepath* and (+++).

Suppose $g = l1 +++ l2$, usually, the key behind the success of this framework is whether we can prove $z \notin \text{path-image } l1$, $z \notin \text{path-image } l2$ and calculate *cindex-pathE* $l1 z$ and *cindex-pathE* $l2 z$.

```

method eval-winding =
  ((rule-tac winding-eq-intro;
    rep-subst
  )
  , auto simp only:winding-simps del:notI intro!:winding-intros
  , tactic (distinct-subgoals-tac))

```

```

end

```

7 Some examples of applying the method winding_eval

```

theory Winding-Number-Eval-Examples imports Winding-Number-Eval
begin

```

```

lemma example1:
  assumes  $R > 1$ 
  shows winding-number (part-circlepath 0 R 0 pi +++ linepath (-R) R) i = 1
  <proof>

lemma example2:
  assumes  $R > 1$ 
  shows winding-number (part-circlepath 0 R 0 pi +++ linepath (-R) R) (-i) =
  0
  <proof>

lemma example3:
  fixes  $lb\ ub\ z :: \text{complex}$ 
  defines  $rec \equiv \text{linepath } lb\ (\text{Complex } (Re\ ub)\ (Im\ lb))\ +++\ \text{linepath } (\text{Complex } (Re\ ub)\ (Im\ lb))\ ub\ +++\ \text{linepath } ub\ (\text{Complex } (Re\ lb)\ (Im\ ub))\ +++\ \text{linepath } (\text{Complex } (Re\ lb)\ (Im\ ub))\ lb$ 
  assumes order-asms:  $Re\ lb < Re\ z < Re\ ub\ Im\ lb < Im\ z < Im\ ub$ 
  shows winding-number  $rec\ z = 1$ 
  <proof>

end

```

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References

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