

Fundamental Properties of Valuation Theory and Hensel's Lemma

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Abstract

Convergence with respect to a valuation is discussed as convergence of a Cauchy sequence. Cauchy sequences of polynomials are defined. They are used to formalize Hensel's lemma.

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```

theory Valuation1
imports Group–Ring–Module.Algebra9
begin

declare ex-image-cong-iff [simp del]

```

Chapter 1

Preliminaries

1.1 Int and ant (augmented integers)

```
lemma int-less-mono:  $(a::nat) < b \implies int\ a < int\ b$   
apply simp  
done
```

```
lemma zless-trans:  $\llbracket (i::int) < j; j < k \rrbracket \implies i < k$   
apply simp  
done
```

```
lemma zmult-pos-bignumTr0:  $\exists L. \forall m. L < m \longrightarrow z < x + int\ m$   
by (subgoal-tac  $\forall m. (nat((abs\ z) + (abs\ x))) < m \longrightarrow z < x + int\ m$ ,  
    blast, rule allI, rule impI, arith)
```

```
lemma zle-less-trans:  $\llbracket (i::int) \leq j; j < k \rrbracket \implies i < k$   
apply (simp add:less-le)  
done
```

```
lemma zless-le-trans:  $\llbracket (i::int) < j; j \leq k \rrbracket \implies i < k$   
apply (simp add:less-le)  
done
```

```
lemma zmult-pos-bignumTr:  $0 < (a::int) \implies$   
     $\exists l. \forall m. l < m \longrightarrow z < x + (int\ m) * a$   
apply (cut-tac zmult-pos-bignumTr0 [of  $z\ x$ ])  
  apply (erule exE)  
  apply (subgoal-tac  $\forall m. L < m \longrightarrow z < x + int\ m * a$ , blast)  
apply (rule allI, rule impI)  
  apply (drule-tac  $a = m$  in forall-spec, assumption)  
  apply (subgoal-tac  $0 \leq int\ m$ )  
  apply (frule-tac  $a = int\ m$  and  $b = a$  in pos-zmult-pos, assumption)  
  apply (cut-tac order-refl [of  $x$ ])  
  apply (frule-tac  $z' = int\ m$  and  $z = int\ m * a$  in  
    zadd-zle-mono [of  $x\ x$ ], assumption+)
```

apply (*rule-tac* $y = x + \text{int } m$ **and** $z = x + (\text{int } m) * a$ **in**
less-le-trans[*of* z], *assumption*+)
apply *simp*
done

lemma *ale-shift*: $\llbracket (x::\text{ant}) \leq y; y = z \rrbracket \implies x \leq z$
by *simp*

lemma *aneg-na-0*[*simp*]: $a < 0 \implies na \ a = 0$
by (*simp add:na-def*)

lemma *amult-an-an:an* $(m * n) = (an \ m) * (an \ n)$
apply (*simp add:an-def*)
apply (*simp add: of-nat-mult a-z-z*)
done

definition
adiv :: $[ant, ant] \Rightarrow ant$ (**infixl** $\langle adiv \rangle$ 200) **where**
 $x \ adiv \ y = ant \ ((tna \ x) \ div \ (tna \ y))$

definition
amod :: $[ant, ant] \Rightarrow ant$ (**infixl** $\langle amod \rangle$ 200) **where**
 $x \ amod \ y = ant \ ((tna \ x) \ mod \ (tna \ y))$

lemma *apos-amod-conj*: $0 < ant \ b \implies$
 $0 \leq (ant \ a) \ amod \ (ant \ b) \wedge (ant \ a) \ amod \ (ant \ b) < (ant \ b)$
by (*simp add:amod-def tna-ant, simp only:ant-0[THEN sym],*
simp add:aless-zless)

lemma *amod-adiv-equality*:
 $(ant \ a) = (a \ div \ b) *_a (ant \ b) + ant \ (a \ mod \ b)$
apply (*simp add:adiv-def tna-ant a-z-z a-zpz asprod-mult*)
done

lemma *asp-z-Z*: $z *_a ant \ x \in Z_\infty$
by (*simp add:asprod-mult z-in-aug-inf*)

lemma *apos-in-aug-inf*: $0 \leq a \implies a \in Z_\infty$
by (*simp add:aug-inf-def, rule contrapos-pp, simp+,*
cut-tac minf-le-any[of 0], frule ale-antisym[of 0 $-\infty$],
assumption+, simp)

lemma *amult-1-both*: $\llbracket 0 < (w::ant); x * w = 1 \rrbracket \implies x = 1 \wedge w = 1$
apply (*cut-tac mem-ant[of x], cut-tac mem-ant[of w],*
(erule disjE)+, simp,
(frule sym, thin-tac $\infty = 1$, simp only:ant-1[THEN sym],
simp del:ant-1))
apply (*erule disjE, erule exE, simp,*
(frule sym, thin-tac $-\infty = 1$, simp only:ant-1[THEN sym],

```

      simp del:ant-1), simp)
apply (frule sym, thin-tac  $-\infty = 1$ , simp only:ant-1[THEN sym],
      simp del:ant-1)
apply ((erule disjE)+, erule exE, simp,
      frule-tac aless-imp-le[of 0  $-\infty$ ],
      cut-tac minf-le-any[of 0],
      frule ale-antisym[of 0  $-\infty$ ], assumption+,
      simp only:ant-0[THEN sym], simp,
      frule sym, thin-tac  $-\infty = 1$ , simp only:ant-1[THEN sym],
      simp del:ant-1)
apply ((erule disjE)+, (erule exE)+, simp only:ant-1[THEN sym],
      simp del:ant-1 add:a-z-z,
      (cut-tac a = z and b = za in mult.commute, simp,
      cut-tac z = za and z' = z in times-1-both, assumption+),
      simp)
apply (erule exE, simp,
      cut-tac x = z and y = 0 in less-linear, erule disjE, simp,
      frule sym, thin-tac  $-\infty = 1$ , simp only:ant-1[THEN sym],
      simp del:ant-1,
      erule disjE, simp add:ant-0, simp,
      frule sym, thin-tac  $\infty = 1$ , simp only:ant-1[THEN sym],
      simp del:ant-1,
      erule disjE, erule exE, simp,
      frule sym, thin-tac  $\infty = 1$ , simp only:ant-1[THEN sym],
      simp del:ant-1, simp)

```

done

lemma poss-int-neq-0: $0 < (z::int) \implies z \neq 0$
by simp

lemma aadd-neg-negg[simp]: $\llbracket a \leq (0::ant); b < 0 \rrbracket \implies a + b < 0$
apply (frule ale-minus[of a 0], simp,
 frule aless-minus[of b 0], simp)
apply (frule aadd-pos-poss[of $-a -b$], assumption+,
 simp add:aminus-add-distrib[THEN sym, of a b],
 frule aless-minus[of 0 $-(a + b)$], simp add:a-minus-minus)
done

lemma aadd-two-negg[simp]: $\llbracket a < (0::ant); b < 0 \rrbracket \implies a + b < 0$
by auto

lemma amin-aminTr: $(z::ant) \leq z' \implies \text{amin } z \ w \leq \text{amin } z' \ w$
by (simp add:amin-def)

lemma amin-le1: $(z::ant) \leq z' \implies (\text{amin } z \ w) \leq z'$
by (simp add:amin-def, simp add:aneg-le,
 rule impI, frule aless-le-trans[of w z z'],
 assumption+, simp add:aless-imp-le)

lemma *amin-le2*:($z::\text{ant}$) $\leq z' \implies (\text{amin } w \ z) \leq z'$
by (*simp add:amin-def*, *rule impI*,
frule ale-trans[of w z z'], *assumption*+)

lemma *Amin-geTr*:($\forall j \leq n. f \ j \in Z_\infty$) $\wedge$ ($\forall j \leq n. z \leq (f \ j)$) $\longrightarrow$
 $z \leq (\text{Amin } n \ f)$

apply (*induct-tac n*)
apply (*rule impI*, *erule conjE*, *simp*)
apply (*rule impI*, (*erule conjE*)+,
cut-tac Nsetn-sub-mem1 [of n], *simp*,
drule-tac x = Suc n in spec, *simp*,
rule-tac z = z and x = Amin n f and y = f(Suc n) in amin-ge1,
simp+)

done

lemma *Amin-ge*: $\llbracket \forall j \leq n. f \ j \in Z_\infty; \forall j \leq n. z \leq (f \ j) \rrbracket \implies$
 $z \leq (\text{Amin } n \ f)$

by (*simp add:Amin-geTr*)

definition

Abs :: *ant* $\Rightarrow$ *ant* **where**
Abs *z* = (*if* *z* < 0 *then* -*z* *else* *z*)

lemma *Abs-pos*:0 $\leq \text{Abs } z$
by (*simp add:Abs-def*, *rule conjI*, *rule impI*,
cut-tac aless-minus[of z 0], *simp*,
assumption,
rule impI, *simp add:aneg-less[of z 0]*)

lemma *Abs-x-plus-x-pos*:0 $\leq (\text{Abs } x) + x$
apply (*case-tac x < 0*,
simp add:Abs-def, *simp add:aadd-minus-inv*)

apply (*simp add:aneg-less*,
simp add:Abs-def, *simp add:aneg-less[THEN sym, of 0 x]*,
simp add:aneg-less[of x 0], *simp add:aadd-two-pos*)

done

lemma *Abs-ge-self*: $x \leq \text{Abs } x$
apply (*simp add:Abs-def*, *rule impI*,
cut-tac ale-minus[of x 0],
simp add:aminus-0, *simp add:aless-imp-le*)

done

lemma *na-1*:*na* 1 = *Suc* 0
apply (*simp only:ant-1[THEN sym]*, *simp only:na-def*,
simp only:ant-0[THEN sym], *simp only:aless-zless[of 1 0]*,
simp, *subgoal-tac* $\infty \neq 1$, *simp*)
apply (*simp only:ant-1[THEN sym]*, *simp only:tna-ant*,

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      rule not-sym, simp only:ant-1[THEN sym], simp del:ant-1)
done

lemma ant-int:ant (int n) = an n
by (simp add:an-def)

lemma int-nat:0 < z  $\implies$  int (nat z) = z
by arith

lemma int-ex-nat:0 < z  $\implies$   $\exists n.$  int n = z
by (cut-tac int-nat[of z], blast, assumption)

lemma eq-nat-pos-ints:
 $\llbracket \text{nat } (z::\text{int}) = \text{nat } (z'::\text{int}); 0 \leq z; 0 \leq z' \rrbracket \implies z = z'$ 
by simp

lemma a-p1-gt[simp]: $\llbracket a \neq \infty; a \neq -\infty \rrbracket \implies a < a + 1$ 
apply (cut-tac aadd-poss-less[of a 1],
      simp add:aadd-commute, assumption+)
apply (cut-tac zposs-aposss[of 1], simp)
done

lemma gt-na-poss:(na a) < m  $\implies$  0 < m
apply (simp add:na-def)
done

lemma azmult-less: $\llbracket a \neq \infty; \text{na } a < m; 0 < x \rrbracket$ 
 $\implies a < \text{int } m *_a x$ 
apply (cut-tac mem-ant[of a])
apply (erule disjE)
apply (case-tac x =  $\infty$ ) apply simp
apply (subst less-le[of  $-\infty \infty$ ]) apply simp
apply (frule aless-imp-le[of 0 x], frule apos-neq-minf[of x])
apply (cut-tac mem-ant[of x], simp, erule exE, simp)
apply (simp add:asprod-amult a-z-z)
apply (simp, erule exE, simp)

apply (frule-tac a = ant z in gt-na-poss[of - m])
apply (case-tac x =  $\infty$ , simp)
apply (frule aless-imp-le[of 0 x])
apply (frule apos-neq-minf[of x])
apply (cut-tac mem-ant[of x], simp, erule exE,
      simp add:asprod-amult a-z-z)
apply (subst aless-zless)
apply (cut-tac a = ant z in gt-na-poss[of - m], assumption)
apply (smt a0-less-int-conv apos-na-poss int-less-mono int-nat na-def of-nat-0-le-iff
pos-zmult-pos tna-ant z-neq-inf)
done

```


lemma *zmult-gt-one*: $\llbracket 2 \leq m; 0 < xa \rrbracket \implies 1 < \text{int } m * xa$
by (*metis ge2-zmult-pos mult.commute*)

lemma *zmult-pos*: $\llbracket 0 < m; 0 < (a::\text{int}) \rrbracket \implies 0 < (\text{int } m) * a$
by (*frule zmult-zless-mono2[of 0 a int m], simp, simp*)

lemma *ant-int-na*: $\llbracket 0 \leq a; a \neq \infty \rrbracket \implies \text{ant } (\text{int } (na\ a)) = a$
by (*frule an-na[of a], assumption, simp add:an-def*)

lemma *zpos-nat*: $0 \leq (z::\text{int}) \implies \exists n. z = \text{int } n$
apply (*subgoal-tac z = int (nat z)*)
apply *blast apply simp*
done

1.2 nsets

lemma *nsetTr1*: $\llbracket j \in \text{nset } a\ b; j \neq a \rrbracket \implies j \in \text{nset } (\text{Suc } a)\ b$
apply (*simp add:nset-def*)
done

lemma *nsetTr2*: $j \in \text{nset } (\text{Suc } a)\ (\text{Suc } b) \implies j - \text{Suc } 0 \in \text{nset } a\ b$
apply (*simp add:nset-def, erule conjE,*
simp add:skip-im-Tr4[of j b])
done

lemma *nsetTr3*: $\llbracket j \neq \text{Suc } (\text{Suc } 0); j - \text{Suc } 0 \in \text{nset } (\text{Suc } 0)\ (\text{Suc } n) \rrbracket$
 $\implies \text{Suc } 0 < j - \text{Suc } 0$
apply (*simp add:nset-def, erule conjE, subgoal-tac j \neq 0,*
rule contrapos-pp, simp+)
done

lemma *Suc-leD1*: $\text{Suc } m \leq n \implies m < n$
apply (*insert lessI[of m],*
rule less-le-trans[of m Suc m n], assumption+)
done

lemma *leI1*: $n < m \implies \neg ((m::\text{nat}) \leq n)$
apply (*rule contrapos-pp, simp+*)
done

lemma *neg-zle*: $\neg (z::\text{int}) \leq z' \implies z' < z$
apply (*simp add: not-le*)
done

lemma *nset-m-m*: $\text{nset } m\ m = \{m\}$
by (*simp add:nset-def,*
rule equalityI, rule subsetI, simp,
rule subsetI, simp)

lemma *nset-Tr51*: $\llbracket j \in \text{nset } (\text{Suc } 0) (\text{Suc } (\text{Suc } n)); j \neq \text{Suc } 0 \rrbracket$
 $\implies j - \text{Suc } 0 \in \text{nset } (\text{Suc } 0) (\text{Suc } n)$
apply (*simp add:nset-def*, (*erule conjE*)+,
frule-tac m = j and n = Suc (Suc n) and l = Suc 0 in diff-le-mono,
simp)
done

lemma *nset-Tr52*: $\llbracket j \neq \text{Suc } (\text{Suc } 0); \text{Suc } 0 \leq j - \text{Suc } 0 \rrbracket$
 $\implies \neg j - \text{Suc } 0 \leq \text{Suc } 0$
by *auto*

lemma *nset-Suc:nset* (*Suc 0*) (*Suc (Suc n)*) =
nset (Suc 0) (Suc n) $\cup$ {Suc (Suc n)}
by (*auto simp add:nset-def*)

lemma *AinequalityTr0*: $x \neq -\infty \implies \exists L. (\forall N. L < N \longrightarrow$
 $(\text{an } m) < (x + \text{an } N))$
apply (*case-tac x = ∞ , simp add:an-def*)
apply (*cut-tac mem-ant[of x], simp, erule exE, simp add:an-def a-zpz*,
simp add:aless-zless,
cut-tac x = z in zmult-pos-bignumTr0[of int m], simp)
done

lemma *AinequalityTr*: $\llbracket 0 < b \wedge b \neq \infty; x \neq -\infty \rrbracket \implies \exists L. (\forall N. L < N \longrightarrow$
 $(\text{an } m) < (x + (\text{int } N) *_a b))$
apply (*frule-tac AinequalityTr0[of x m]*,
erule exE,
*subgoal-tac $\forall N. L < N \longrightarrow \text{an } m < x + (\text{int } N) *_a b$* ,
blast, rule allI, rule impI)
apply (*drule-tac a = N in forall-spec, assumption*,
erule conjE,
cut-tac N = N in asprod-ge[of b], assumption,
thin-tac x $\neq -\infty$, thin-tac b $\neq \infty$, thin-tac an m < x + an N,
simp)
apply (*frule-tac x = an N and y = int N *_a b and z = x in aadd-le-mono*,
simp only:aadd-commute[of - x])
done

lemma *two-inequalities*: $\llbracket \forall (n::\text{nat}). x < n \longrightarrow P \ n; \forall (n::\text{nat}). y < n \longrightarrow Q \ n \rrbracket$
 $\implies \forall n. (\max x \ y) < n \longrightarrow (P \ n) \wedge (Q \ n)$
by *auto*

lemma *multi-inequalityTr0*: $(\forall j \leq (n::\text{nat}). (x \ j) \neq -\infty) \longrightarrow$
 $(\exists L. (\forall N. L < N \longrightarrow (\forall l \leq n. (\text{an } m) < (x \ l) + (\text{an } N))))$
apply (*induct-tac n*)
apply (*rule impI, simp*)
apply (*rule AinequalityTr0[of x 0 m], assumption*)
apply (*rule impI*)

```

apply (subgoal-tac  $\forall l. l \leq n \longrightarrow l \leq (\text{Suc } n)$ , simp)
apply (erule exE)
apply (frule-tac  $a = \text{Suc } n$  in forall-spec, simp)

apply (frule-tac  $x = x (\text{Suc } n)$  in AinequalityTr0[of - m])
apply (erule exE)
apply (subgoal-tac  $\forall N. (\max L La) < N \longrightarrow$ 
   $(\forall l \leq (\text{Suc } n). \text{an } m < x l + \text{an } N)$ , blast)
apply (rule allI, rule impI, rule allI, rule impI)
apply (rotate-tac 1)
apply (case-tac  $l = \text{Suc } n$ , simp,
  drule-tac  $m = l$  and  $n = \text{Suc } n$  in noteq-le-less, assumption+,
  drule-tac  $x = l$  and  $n = \text{Suc } n$  in less-le-diff, simp,
  simp)
done

lemma multi-inequalityTr1:  $\llbracket \forall j \leq (n::\text{nat}). (x j) \neq -\infty \rrbracket \implies$ 
 $\exists L. (\forall N. L < N \longrightarrow (\forall l \leq n. (\text{an } m) < (x l) + (\text{an } N)))$ 
by (simp add: multi-inequalityTr0)

lemma gcoeff-multi-inequality:  $\llbracket \forall N. 0 < N \longrightarrow (\forall j \leq (n::\text{nat}). (x j) \neq -\infty \wedge$ 
 $0 < (b N j) \wedge (b N j) \neq \infty) \rrbracket \implies$ 
 $\exists L. (\forall N. L < N \longrightarrow (\forall l \leq n. (\text{an } m) < (x l) + (\text{int } N) *_a (b N l)))$ 
apply (subgoal-tac  $\forall j \leq n. x j \neq -\infty$ )
apply (frule multi-inequalityTr1[of n x m])
apply (erule exE)
apply (subgoal-tac  $\forall N. L < N \longrightarrow$ 
 $(\forall l \leq n. \text{an } m < x l + (\text{int } N) *_a (b N l)))$ 
apply blast

apply (rule allI, rule impI, rule allI, rule impI,
  drule-tac  $a = N$  in forall-spec, simp,
  drule-tac  $a = l$  in forall-spec, assumption,
  drule-tac  $a = N$  in forall-spec, assumption,
  drule-tac  $a = l$  in forall-spec, assumption,
  drule-tac  $a = l$  in forall-spec, assumption)
apply (cut-tac  $b = b N l$  and  $N = N$  in asprod-ge, simp, simp,
  (erule conjE)+, simp, thin-tac  $x l \neq -\infty$ , thin-tac  $b N l \neq \infty$ )
apply (frule-tac  $x = \text{an } N$  and  $y = \text{int } N *_a b N l$  and  $z = x l$  in
  aadd-le-mono, simp add: aadd-commute,
  rule allI, rule impI,
  cut-tac lessI[of (0::nat)],
  drule-tac  $a = \text{Suc } 0$  in forall-spec, assumption)
apply simp
done

primrec m-max ::  $[\text{nat}, \text{nat} \Rightarrow \text{nat}] \Rightarrow \text{nat}$ 
where
  m-max-0: m-max 0 f = f 0

```

| *m-max-Suc*: $m\text{-max } (Suc\ n)\ f = \max\ (m\text{-max } n\ f)\ (f\ (Suc\ n))$

lemma *m-maxTr*: $\forall l \leq n. (f\ l) \leq m\text{-max } n\ f$
apply (*induct-tac* *n*)
apply *simp*

apply (*rule allI*, *rule impI*)
apply *simp*
apply (*case-tac* $l = Suc\ n$, *simp*)
apply (*cut-tac* $m = l$ **and** $n = Suc\ n$ **in** *noteq-le-less*, *assumption*+,
 $thin\text{-tac } l \leq Suc\ n$, $thin\text{-tac } l \neq Suc\ n$,
 $frule\text{-tac } x = l$ **and** $n = Suc\ n$ **in** *less-le-diff*,
 $thin\text{-tac } l < Suc\ n$, *simp*)
apply (*drule-tac* $a = l$ **in** *forall-spec*, *assumption*)
apply *simp*
done

lemma *m-max-gt*: $l \leq n \implies (f\ l) \leq m\text{-max } n\ f$
apply (*simp add*:*m-maxTr*)
done

lemma *ASum-zero*: $(\forall j \leq n. f\ j \in Z_\infty) \wedge (\forall l \leq n. f\ l = 0) \longrightarrow A\text{Sum } f\ n = 0$
apply (*induct-tac* *n*)
apply (*rule impI*, *erule conjE*, *simp*)
apply (*rule impI*)
apply (*subgoal-tac* $(\forall j \leq n. f\ j \in Z_\infty) \wedge (\forall l \leq n. f\ l = 0)$, *simp*)
apply (*simp add*:*aadd-0-l*, *erule conjE*,
 $thin\text{-tac } (\forall j \leq n. f\ j \in Z_\infty) \wedge (\forall l \leq n. f\ l = 0) \longrightarrow A\text{Sum } f\ n = 0$)
apply (*rule conjI*)
apply (*rule allI*, *rule impI*,
 $drule\text{-tac } a = j$ **in** *forall-spec*, *simp*, *assumption*+)
apply ($thin\text{-tac } \forall j \leq Suc\ n. f\ j \in Z_\infty$)
apply (*rule allI*, *rule impI*,
 $drule\text{-tac } a = l$ **in** *forall-spec*, *simp*+)
done

lemma *eSum-singleTr*: $(\forall j \leq n. f\ j \in Z_\infty) \wedge (j \leq n \wedge (\forall l \in \{h. h \leq n\} - \{j\}. f\ l = 0)) \longrightarrow A\text{Sum } f\ n = f\ j$
apply (*induct-tac* *n*)
apply (*simp*, *rule impI*, (*erule conjE*)**+**)
apply (*case-tac* $j \leq n$)
apply *simp*
apply (*simp add*:*aadd-0-r*)
apply *simp*
apply (*simp add*:*nat-not-le-less*[*of j*])
apply ($frule\text{-tac } m = n$ **and** $n = j$ **in** *Suc-leI*)
apply ($frule\text{-tac } m = j$ **and** $n = Suc\ n$ **in** *le-antisym*, *assumption*+, *simp*)

apply (*cut-tac* $n = n$ **in** *ASum-zero* [*of* - *f*])
apply (*subgoal-tac* $(\forall j \leq n. f\ j \in Z_\infty) \wedge (\forall l \leq n. f\ l = 0)$)
apply (*thin-tac* $\forall j \leq \text{Suc } n. f\ j \in Z_\infty$,
thin-tac $\forall l \in \{h. h \leq \text{Suc } n\} - \{\text{Suc } n\}. f\ l = 0$, *simp only:mp*)
apply (*simp add:aadd-0-l*)

apply (*thin-tac* $(\forall j \leq n. f\ j \in Z_\infty) \wedge (\forall l \leq n. f\ l = 0) \longrightarrow \text{ASum } f\ n = 0$)
apply (*rule conjI*,
thin-tac $\forall l \in \{h. h \leq \text{Suc } n\} - \{\text{Suc } n\}. f\ l = 0$, *simp*)
apply (*thin-tac* $\forall j \leq \text{Suc } n. f\ j \in Z_\infty$, *simp*)
done

lemma *eSum-single*: $\llbracket \forall j \leq n. f\ j \in Z_\infty ; j \leq n ; \forall l \in \{h. h \leq n\} - \{j\}. f\ l = 0 \rrbracket$
 $\implies \text{ASum } f\ n = f\ j$
apply (*simp add:eSum-singleTr*)
done

lemma *ASum-eqTr*: $(\forall j \leq n. f\ j \in Z_\infty) \wedge (\forall j \leq n. g\ j \in Z_\infty) \wedge$
 $(\forall j \leq n. f\ j = g\ j) \longrightarrow \text{ASum } f\ n = \text{ASum } g\ n$
apply (*induct-tac* *n*)
apply (*rule impI*, *simp*)

apply (*rule impI*, (*erule conjE*)+)
apply *simp*
done

lemma *ASum-eq*: $\llbracket \forall j \leq n. f\ j \in Z_\infty ; \forall j \leq n. g\ j \in Z_\infty ; \forall j \leq n. f\ j = g\ j \rrbracket \implies$
 $\text{ASum } f\ n = \text{ASum } g\ n$
by (*cut-tac* *ASum-eqTr*[*of* *n f g*], *simp*)

definition

Kronecker-delta :: $[\text{nat}, \text{nat}] \Rightarrow \text{ant}$
 $(\langle \delta _ _ \rangle \ [70,71] \ 70)$ **where**
 $\delta_i\ j = (\text{if } i = j \text{ then } 1 \text{ else } 0)$

definition

K-gamma :: $[\text{nat}, \text{nat}] \Rightarrow \text{int}$
 $(\langle \gamma _ _ \rangle \ [70,71] \ 70)$ **where**
 $\gamma_i\ j = (\text{if } i = j \text{ then } 0 \text{ else } 1)$

abbreviation

TRANSPOS $(\langle \tau _ _ \rangle \ [90,91] \ 90)$ **where**
 $\tau_i\ j == \text{transpos } i\ j$

lemma *Kdelta-in-Zinf*: $\llbracket j \leq (\text{Suc } n) ; k \leq (\text{Suc } n) \rrbracket \implies$
 $z *_{\text{a}} (\delta_j\ k) \in Z_\infty$
apply (*simp add:Kronecker-delta-def*)
apply (*simp add:z-in-aug-inf Zero-in-aug-inf*)

apply (*simp add:asprod-n-0 Zero-in-aug-inf*)
done

lemma *Kdelta-in-Zinf1*: $\llbracket j \leq n; k \leq n \rrbracket \implies \delta_j k \in Z_\infty$
apply (*simp add:Kronecker-delta-def*)
apply (*simp add:z-in-aug-inf Zero-in-aug-inf*)
apply (*rule impI*)
apply (*simp only:ant-1[THEN sym], simp del:ant-1 add:z-in-aug-inf*)
done

primrec *m-zmax* :: $[nat, nat \Rightarrow int] \Rightarrow int$
where
m-zmax-0: *m-zmax* 0 *f* = *f* 0
| *m-zmax-Suc*: *m-zmax* (*Suc* *n*) *f* = *zmax* (*m-zmax* *n* *f*) (*f* (*Suc* *n*))

lemma *m-zmax-gt-eachTr*:
 $(\forall j \leq n. f j \in Zset) \longrightarrow (\forall j \leq n. (f j) \leq m\text{-}zmax\ n\ f)$
apply (*induct-tac n*)
apply (*rule impI, rule allI, rule impI, simp*)
apply (*rule impI*)
apply *simp*
apply (*rule allI, rule impI*)
apply (*case-tac j = Suc n, simp*)
apply (*simp add:zmax-def*)
apply (*drule-tac m = j and n = Suc n in noteq-le-less, assumption,*
drule-tac x = j and n = Suc n in less-le-diff, simp)
apply (*drule-tac a = j in forall-spec, assumption*)
apply (*simp add:zmax-def*)
done

lemma *m-zmax-gt-each*: $(\forall j \leq n. f j \in Zset) \implies (\forall j \leq n. (f j) \leq m\text{-}zmax\ n\ f)$
apply (*simp add:m-zmax-gt-eachTr*)
done

lemma *n-notin-Nset-pred*: $0 < n \implies \neg n \leq (n - Suc\ 0)$
apply *simp*
done

lemma *Nset-preTr*: $\llbracket 0 < n; j \leq (n - Suc\ 0) \rrbracket \implies j \leq n$
apply *simp*
done

lemma *Nset-preTr1*: $\llbracket 0 < n; j \leq (n - Suc\ 0) \rrbracket \implies j \neq n$
apply *simp*
done

lemma *transpos-noteqTr*: $\llbracket 0 < n; k \leq (n - Suc\ 0); j \leq n; j \neq n \rrbracket$
 $\implies j \neq (\tau_j\ n)\ k$
apply (*rule contrapos-pp, simp+*)

```
apply (simp add:transpos-def)
apply (case-tac k = j, simp, simp)
apply (case-tac k = n, simp)
apply (simp add:n-notin-Nset-pred)
done
```

Chapter 2

Elementary properties of a valuation

2.1 Definition of a valuation

definition

$valuation :: [('b, 'm) \text{ Ring-scheme}, 'b \Rightarrow ant] \Rightarrow bool$ **where**
 $valuation\ K\ v \longleftrightarrow$
 $v \in extensional\ (carrier\ K) \wedge$
 $v \in carrier\ K \rightarrow Z_\infty \wedge$
 $v\ (\mathbf{0}_K) = \infty \wedge (\forall x \in (carrier\ K) - \{\mathbf{0}_K\}. v\ x \neq \infty) \wedge$
 $(\forall x \in (carrier\ K). \forall y \in (carrier\ K). v\ (x \cdot_r K\ y) = (v\ x) + (v\ y)) \wedge$
 $(\forall x \in (carrier\ K). 0 \leq (v\ x) \longrightarrow 0 \leq (v\ (I_r K \pm_K x))) \wedge$
 $(\exists x. x \in carrier\ K \wedge (v\ x) \neq \infty \wedge (v\ x) \neq 0)$

lemma (in *Corps*) *invf-closed*: $x \in carrier\ K - \{\mathbf{0}\} \implies x^{-K} \in carrier\ K$
by (*cut-tac invf-closed1*[*of x*], *simp*, *assumption*)

lemma (in *Corps*) *valuation-map*: $valuation\ K\ v \implies v \in carrier\ K \rightarrow Z_\infty$
by (*simp add: valuation-def*)

lemma (in *Corps*) *value-in-aug-inf*: $\llbracket valuation\ K\ v; x \in carrier\ K \rrbracket \implies$
 $v\ x \in Z_\infty$
by (*simp add: valuation-def*, (*erule conjE*) $+$, *simp add: funcset-mem*)

lemma (in *Corps*) *value-of-zero*: $valuation\ K\ v \implies v\ (\mathbf{0}) = \infty$
by (*simp add: valuation-def*)

lemma (in *Corps*) *val-nonzero-noninf*: $\llbracket valuation\ K\ v; x \in carrier\ K; x \neq \mathbf{0} \rrbracket$
 $\implies (v\ x) \neq \infty$
by (*simp add: valuation-def*)

lemma (in *Corps*) *value-inf-zero*: $\llbracket valuation\ K\ v; x \in carrier\ K; v\ x = \infty \rrbracket$
 $\implies x = \mathbf{0}$
by (*rule contrapos-pp*, *simp* $+$,

$\text{frule val-nonzero-noninf[of } v \ x], \text{ assumption+}, \text{ simp})$

lemma (in Corps) val-nonzero-z: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; x \neq \mathbf{0} \rrbracket \implies$
 $\exists z. (v \ x) = \text{ant } z$

by (frule value-in-aug-inf[of $v \ x$], assumption+,
frule val-nonzero-noninf[of $v \ x$], assumption+,
cut-tac mem-ant[of $v \ x$], simp add:aug-inf-def)

lemma (in Corps) val-nonzero-z-unique: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; x \neq \mathbf{0} \rrbracket$
 $\implies \exists! z. (v \ x) = \text{ant } z$

by (rule ex-ex1I, simp add:val-nonzero-z, simp)

lemma (in Corps) value-noninf-nonzero: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; v \ x \neq \infty \rrbracket$
 $\implies x \neq \mathbf{0}$

by (rule contrapos-pp, simp+, simp add:value-of-zero)

lemma (in Corps) val1-neq-0: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; v \ x = 1 \rrbracket \implies$
 $x \neq \mathbf{0}$

apply (rule contrapos-pp, simp+, simp add:value-of-zero)
apply (simp only:ant-1[THEN sym], cut-tac z-neq-inf[THEN not-sym, of 1], simp)
done

lemma (in Corps) val-Zmin-sym: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; y \in \text{carrier } K \rrbracket$
 $\implies \text{amin } (v \ x) \ (v \ y) = \text{amin } (v \ y) \ (v \ x)$

by (simp add:amin-commute)

lemma (in Corps) val-t2p: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; y \in \text{carrier } K \rrbracket$
 $\implies v \ (x \cdot_r y) = v \ x + v \ y$

by (simp add:valuation-def)

lemma (in Corps) val-axiom4: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; 0 \leq v \ x \rrbracket \implies$
 $0 \leq v \ (1_r \pm x)$

by (simp add:valuation-def)

lemma (in Corps) val-axiom5: $\text{valuation } K \ v \implies$
 $\exists x. x \in \text{carrier } K \wedge v \ x \neq \infty \wedge v \ x \neq 0$

by (simp add:valuation-def)

lemma (in Corps) val-field-nonzero: $\text{valuation } K \ v \implies \text{carrier } K \neq \{\mathbf{0}\}$

by (rule contrapos-pp, simp+,
frule val-axiom5[of v],
erule exE, (erule conjE)+, simp add:value-of-zero)

lemma (in Corps) val-field-1-neq-0: $\text{valuation } K \ v \implies 1_r \neq \mathbf{0}$

apply (rule contrapos-pp, simp+)
apply (frule val-axiom5[of v])
apply (erule exE, (erule conjE)+)
apply (cut-tac field-is-ring,
frule-tac $t = x$ in Ring.ring-l-one[THEN sym, of K], assumption+,

simp add:Ring.ring-times-0-x, simp add:value-of-zero)
done

lemma (in Corps) *value-of-one:valuation* $K\ v \implies v\ (1_r) = 0$
apply (*cut-tac field-is-ring, frule Ring.ring-one*[of K])
apply (*frule val-t2p*[of $v\ 1_r\ 1_r$], *assumption+*,
simp add:Ring.ring-l-one, frule val-field-1-neq-0[of v],
frule val-nonzero-z[of $v\ 1_r$], *assumption+*,
erule exE, simp add:a-zpz)
done

lemma (in Corps) *has-val-one-neq-zero:valuation* $K\ v \implies 1_r \neq 0$
by (*frule value-of-one*[of v],
rule contrapos-pp, simp+, simp add:value-of-zero)

lemma (in Corps) *val-minus-one:valuation* $K\ v \implies v\ (-_a\ 1_r) = 0$
apply (*cut-tac field-is-ring, frule Ring.ring-one*[of K],
frule Ring.ring-is-ag[of K],
frule val-field-1-neq-0[of v],
frule aGroup.ag-inv-inj[of $K\ 1_r\ 0$], *assumption+*,
simp add:Ring.ring-zero, assumption)
apply (*frule val-nonzero-z*[of $v\ -_a\ 1_r$],
rule aGroup.ag-mOp-closed, assumption+, simp add:aGroup.ag-inv-zero,
erule exE, frule val-t2p [THEN sym, of $v\ -_a\ 1_r\ -_a\ 1_r$])
apply (*simp add:aGroup.ag-mOp-closed*[of $K\ 1_r$],
simp add:aGroup.ag-mOp-closed[of $K\ 1_r$],
frule Ring.ring-inv1-3[THEN sym, of $K\ 1_r\ 1_r$], assumption+,
simp add:Ring.ring-l-one, simp add:value-of-one a-zpz)
done

lemma (in Corps) *val-minus-eq:valuation* $K\ v; x \in \text{carrier } K \implies$
 $v\ (-_a\ x) = v\ x$
apply (*cut-tac field-is-ring,*
simp add:Ring.ring-times-minusl[of $K\ x$],
subst val-t2p[of v], *assumption+*,
frule Ring.ring-is-ag[of K], *rule aGroup.ag-mOp-closed, assumption+,*
simp add:Ring.ring-one, assumption, simp add:val-minus-one,
simp add:aadd-0-l)
done

lemma (in Corps) *value-of-inv:valuation* $K\ v; x \in \text{carrier } K; x \neq 0 \implies$
 $v\ (x^{-K}) = -\ (v\ x)$
apply (*cut-tac invf-inv*[of x], *erule conjE,*
frule val-t2p[of $v\ x^{-K}\ x$], *assumption+*,
simp+, simp add:value-of-one, simp add:a-inv)
apply simp
done

lemma (in Corps) *val-exp-ring:valuation* $K\ v; x \in \text{carrier } K; x \neq 0$

$\implies (\text{int } n) *_{\mathbf{a}} (v \ x) = v \ (x^{\sim K} n)$
apply (*cut-tac field-is-ring*,
 $\text{induct-tac } n, \text{ simp add: Ring.ring-r-one, simp add: value-of-one}$)
apply (*drule sym, simp*)
apply (*subst val-t2p[$\text{of } v - x$], assumption+,*
 $\text{rule Ring.npClose, assumption+,}$
 $\text{frule val-nonzero-z[$\text{of } v \ x$], assumption+,}$
 $\text{erule exE, simp add: asprod-mult a-zpz,}$
 $\text{simp add: distrib-right}$)
done

exponent in a field

lemma (*in Corps*) $\text{val-exp} : \llbracket \text{valuation } K \ v; x \in \text{carrier } K; x \neq \mathbf{0} \rrbracket \implies$
 $z *_{\mathbf{a}} (v \ x) = v \ (x_K^z)$
apply (*simp add: npowf-def*)
apply (*case-tac $0 \leq z$,*
 $\text{simp, frule val-exp-ring [of } v \ x \ \text{nat } z], \text{ assumption+,}$
 simp, simp)
apply (*simp add: zle,*
 $\text{cut-tac invf-closed1 [of } x], \text{ simp,}$
 $\text{cut-tac val-exp-ring [THEN sym, of } v \ x^{-K} \ \text{nat } (-z)], \text{ simp,}$
 $\text{thin-tac } v \ (x^{-K} \sim K \ (\text{nat } (-z))) = (-z) *_{\mathbf{a}} v \ (x^{-K}), \text{erule conjE}$)
apply (*subst value-of-inv[$\text{of } v \ x$], assumption+*)
apply (*frule val-nonzero-z[$\text{of } v \ x$], assumption+,erule exE, simp,*
 $\text{simp add: asprod-mult aminus, simp+}$)
done

lemma (*in Corps*) $\text{value-zero-nonzero} : \llbracket \text{valuation } K \ v; x \in \text{carrier } K; v \ x = 0 \rrbracket$
 $\implies x \neq \mathbf{0}$
by (*frule value-noninf-nonzero[$\text{of } v \ x$], assumption+, simp,*
 assumption)

lemma (*in Corps*) $\text{v-ale-diff} : \llbracket \text{valuation } K \ v; x \in \text{carrier } K; y \in \text{carrier } K;$
 $x \neq \mathbf{0}; v \ x \leq v \ y \rrbracket \implies 0 \leq v(y \cdot_r x^{-K})$
apply (*frule value-in-aug-inf[$\text{of } v \ x$], simp+,*
 $\text{frule value-in-aug-inf[$\text{of } v \ y$], simp+,}$
 $\text{frule val-nonzero-z[$\text{of } v \ x$], assumption+,}$
 erule exE)
apply (*cut-tac invf-closed[$\text{of } x$], simp+,*
 $\text{simp add: val-t2p,}$
 $\text{simp add: value-of-inv[$\text{of } v \ x$],}$
 $\text{frule-tac } x = \text{ant } z \ \text{in ale-diff-pos[$\text{of } -v \ y$],}$
 $\text{simp add: diff-ant-def}$)
apply *simp*
done

lemma (*in Corps*) $\text{amin-le-plusTr} : \llbracket \text{valuation } K \ v; x \in \text{carrier } K; y \in \text{carrier } K;$
 $v \ x \neq \infty; v \ y \neq \infty; v \ x \leq v \ y \rrbracket \implies \text{amin } (v \ x) \ (v \ y) \leq v \ (x \pm y)$
apply (*cut-tac field-is-ring, frule Ring.ring-is-ag,*

```

    frule value-noninf-nonzero[of v x], assumption+,
    frule v-ale-diff[of v x y], assumption+,
    cut-tac invf-closed1[of x],
    frule Ring.ring-tOp-closed[of K y x- K], assumption+, simp,
    frule Ring.ring-one[of K],
    frule aGroup.ag-pOp-closed[of K 1r y ·r x- K], assumption+,
    frule val-axiom4[of v y ·r (x- K)], assumption+)
apply (frule aadd-le-mono[of 0 v (1r ± y ·r x- K) v x],
    simp add:aadd-0-l, simp add:aadd-commute[of - v x],
    simp add:val-t2p[THEN sym, of v x],
    simp add:Ring.ring-distrib1 Ring.ring-r-one,
    simp add:Ring.ring-tOp-commute[of K x],
    simp add:Ring.ring-tOp-assoc, simp add:linvf,
    simp add:Ring.ring-r-one,
    cut-tac amin-le-l[of v x v y],
    rule ale-trans[of amin (v x) (v y) v x v (x ± y)], assumption+)
apply simp
done

lemma (in Corps) amin-le-plus:⟦valuation K v; x ∈ carrier K; y ∈ carrier K⟧
    ⟹ (amin (v x) (v y)) ≤ (v (x ± y))
apply (cut-tac field-is-ring, frule Ring.ring-is-ag)
apply (case-tac v x = ∞ ∨ v y = ∞)
apply (erule disjE, simp,
    frule value-inf-zero[of v x], assumption+,
    simp add:aGroup.ag-l-zero amin-def,
    frule value-inf-zero[of v y], assumption+,
    simp add:aGroup.ag-r-zero amin-def, simp, erule conjE)
apply (cut-tac z = v x and w = v y in ale-linear,
    erule disjE, simp add:amin-le-plusTr,
    frule-tac amin-le-plusTr[of v y x], assumption+,
    simp add:aGroup.ag-pOp-commute amin-commute)
done

lemma (in Corps) value-less-eq:⟦valuation K v; x ∈ carrier K; y ∈ carrier K;
    (v x) < (v y)⟧ ⟹ (v x) = (v (x ± y))
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    frule amin-le-plus[of v x y], assumption+,
    frule aless-imp-le[of v x v y],
    simp add: amin-def)
apply (frule amin-le-plus[of v x ± y -a y],
    rule aGroup.ag-pOp-closed, assumption+,
    rule aGroup.ag-mOp-closed, assumption+,
    simp add:val-minus-eq,
    frule aGroup.ag-mOp-closed[of K y], assumption+,
    simp add:aGroup.ag-pOp-assoc[of K x y],
    simp add:aGroup.ag-r-inv1, simp add:aGroup.ag-r-zero,
    simp add:amin-def)
apply (case-tac ¬ (v (x ±K y) ≤ (v y)), simp+)

```

done

lemma (in Corps) value-less-eq1: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; y \in \text{carrier } K; (v \ x) < (v \ y) \rrbracket \implies v \ x = v \ (y \pm x)$
apply (cut-tac field-is-ring,
frule Ring.ring-is-ag[of K],
frule value-less-eq[of v x y], assumption+)
apply (subst aGroup.ag-pOp-commute, assumption+)
done

lemma (in Corps) val-1px: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; 0 \leq (v \ (1_r \pm x)) \rrbracket \implies 0 \leq (v \ x)$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
frule Ring.ring-one[of K])
apply (rule contrapos-pp, simp+,
case-tac x = 0_K,
simp add: aGroup.ag-r-zero, simp add: value-of-zero,
simp add: aneg-le[of 0 v x],
frule value-less-eq[of v x 1_r], assumption+,
simp add: value-of-one)
apply (drule sym,
simp add: aGroup.ag-pOp-commute[of K x])
done

lemma (in Corps) val-1mx: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; 0 \leq (v \ (1_r \pm (-_a \ x))) \rrbracket \implies 0 \leq (v \ x)$
by (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
frule val-1px[of v -_a x],
simp add: aGroup.ag-mOp-closed, assumption, simp add: val-minus-eq)

2.2 The normal valuation of v

definition

$Lv :: [(r, m) \text{ Ring-scheme}, r \Rightarrow \text{ant}] \Rightarrow \text{ant}$ **where**
 $Lv \ K \ v = AMin \ \{x. x \in v \text{ 'carrier } K \wedge 0 < x\}$

definition

$n\text{-val} :: [(r, m) \text{ Ring-scheme}, r \Rightarrow \text{ant}] \Rightarrow (r \Rightarrow \text{ant})$ **where**
 $n\text{-val} \ K \ v = (\lambda x \in \text{carrier } K. (THE \ l. (l * (Lv \ K \ v)) = v \ x))$

definition

$Pg :: [(r, m) \text{ Ring-scheme}, r \Rightarrow \text{ant}] \Rightarrow r$ **where**
 $Pg \ K \ v = (SOME \ x. x \in \text{carrier } K - \{0_K\} \wedge v \ x = Lv \ K \ v)$

lemma (in Corps) vals-pos-nonempty: $\text{valuation } K \ v \implies \{x. x \in v \text{ 'carrier } K \wedge 0 < x\} \neq \{\}$
using val-axiom5[of v] value-noninf-nonzero[of v] value-of-inv[THEN sym, of v]
by (auto simp: ex-image-cong-iff) (metis Ring.ring-is-ag aGroup.ag-mOp-closed)

aGroup.ag-pOp-closed aGroup.ag-r-inv1 f-is-ring zero-lt-inf)

lemma (in Corps) *vals-pos-LBset:valuation K v $\implies$*
 $\{x. x \in v \text{ ' carrier } K \wedge 0 < x\} \subseteq \text{LBset } 1$
by (rule subsetI, simp add:LBset-def, erule conjE,
rule-tac $x = x$ in gt-a0-ge-1, assumption)

lemma (in Corps) *Lv-pos:valuation K v $\implies 0 < Lv K v$*
apply (simp add:Lv-def,
frule vals-pos-nonempty[of v],
frule vals-pos-LBset[of v],
simp only:ant-1[THEN sym],
frule AMin[of $\{x. x \in v \text{ ' carrier } K \wedge 0 < x\} 1$], assumption+,
erule conjE)
apply simp
done

lemma (in Corps) *AMin-z:valuation K v $\implies$*
 $\exists a. \text{AMin } \{x. x \in v \text{ ' carrier } K \wedge 0 < x\} = \text{ant } a$
apply (frule vals-pos-nonempty[of v],
frule vals-pos-LBset[of v],
simp only:ant-1[THEN sym],
frule AMin[of $\{x. x \in v \text{ ' carrier } K \wedge 0 < x\} 1$], assumption+,
erule conjE)
apply (frule val-axiom5[of v],
erule exE, (erule conjE)+,
cut-tac $x = v x$ in aless-linear[of - 0], simp,
erule disjE,
frule-tac $x = x$ in value-noninf-nonzero[of v], assumption+,
frule-tac $x1 = x$ in value-of-inv[THEN sym, of v], assumption+)
apply (frule-tac $x = v x$ in aless-minus[of - 0], simp,
cut-tac $x = x$ in invf-closed1, simp, erule conjE,
frule valuation-map[of v],
frule-tac $a = x^{-K}$ in mem-in-image[of v carrier K Z_∞], simp)
apply (drule-tac $a = v (x^{-K})$ in forall-spec, simp,
frule-tac $x = x^{-K}$ in val-nonzero-noninf[of v],
thin-tac $v (x^{-K}) \in v \text{ ' carrier } K$,
thin-tac $\{x \in v \text{ ' carrier } K. 0 < x\} \subseteq \text{LBset } 1$,
thin-tac $\text{AMin } \{x \in v \text{ ' carrier } K. 0 < x\} \in v \text{ ' carrier } K$,
thin-tac $0 < \text{AMin } \{x \in v \text{ ' carrier } K. 0 < x\}$, simp,
thin-tac $v (x^{-K}) \in v \text{ ' carrier } K$,
thin-tac $\{x \in v \text{ ' carrier } K. 0 < x\} \subseteq \text{LBset } 1$,
thin-tac $\text{AMin } \{x \in v \text{ ' carrier } K. 0 < x\} \in v \text{ ' carrier } K$,
thin-tac $0 < \text{AMin } \{x \in v \text{ ' carrier } K. 0 < x\}$, simp)
apply (rule noninf-mem-Z[of AMin $\{x \in v \text{ ' carrier } K. 0 < x\}$],
frule image-sub[of v carrier K Z_∞ carrier K],
rule subset-refl)
apply (rule subsetD[of v ' carrier K Z_∞
AMin $\{x \in v \text{ ' carrier } K. 0 < x\}$], assumption+)

apply *auto*
by (*metis* (*no-types*, *lifting*) *aneg-le aug-inf-noninf-is-z image-eqI value-in-aug-inf z-less-i*)

lemma (**in** *Corps*) *Lv-z:valuation K v $\implies \exists z. Lv K v = ant z$*
by (*simp add:Lv-def, rule AMin-z, assumption+*)

lemma (**in** *Corps*) *AMin-k:valuation K v $\implies$*
 $\exists k \in \text{carrier } K - \{0\}. \text{AMin } \{x. x \in v \text{ ' carrier } K \wedge 0 < x\} = v k$

apply (*frule vals-pos-nonempty[of v],*
frule vals-pos-LBset[of v],
simp only:ant-1[THEN sym],
frule AMin[of {x. x $\in$ v ' carrier K $\wedge$ 0 < x} 1], assumption+,
erule conjE)

apply (*thin-tac $\forall x \in \{x. x \in v \text{ ' carrier } K \wedge 0 < x\}.$*
 $\text{AMin } \{x. x \in v \text{ ' carrier } K \wedge 0 < x\} \leq x$)

apply (*simp add:image-def, erule conjE, erule bexE,*
thin-tac {x. ($\exists xa \in \text{carrier } K. x = v xa$) $\wedge$ 0 < x} $\subseteq$ LBset 1,
thin-tac $\exists x. (\exists xa \in \text{carrier } K. x = v xa) \wedge 0 < x,$
subgoal-tac $x \in \text{carrier } K - \{0\},$ blast,
frule AMin-z[of v], erule exE, simp)

apply (*simp add:image-def,*
thin-tac AMin {x. ($\exists xa \in \text{carrier } K. x = v xa$) $\wedge$ 0 < x} = ant a,
rule contrapos-pp, simp+, frule sym, thin-tac $v(0) = ant a,$
simp add:value-of-zero)

done

lemma (**in** *Corps*) *val-Pg: valuation K v $\implies$*
 $Pg K v \in \text{carrier } K - \{0\} \wedge v (Pg K v) = Lv K v$

apply (*frule AMin-k[of v], unfold Lv-def, unfold Pg-def*)

apply (*rule someI2-ex*)

apply (*erule bexE, drule sym, unfold Lv-def, blast*)

apply *simp*

done

lemma (**in** *Corps*) *amin-generateTr:valuation K v $\implies$*

$\forall w \in \text{carrier } K - \{0\}. \exists z. v w = z *_a \text{AMin } \{x. x \in v \text{ ' carrier } K \wedge 0 < x\}$

apply (*frule vals-pos-nonempty[of v],*
frule vals-pos-LBset[of v],
simp only:ant-1[THEN sym],
frule AMin[of {x. x $\in$ v ' carrier K $\wedge$ 0 < x} 1], assumption+,
frule AMin-z[of v], erule exE, simp,
thin-tac $\exists x. x \in v \text{ ' carrier } K \wedge 0 < x,$
(erule conjE)+, rule ballI, simp, erule conjE,
*frule-tac $x = w$ **in** val-nonzero-noninf[of v], assumption+,*
*frule-tac $x = w$ **in** value-in-aug-inf[of v], assumption+,*
simp add:aug-inf-def,
*cut-tac $a = v w$ **in** mem-ant, simp, erule exE,*

$\text{cut-tac } a = z \text{ and } b = a \text{ in } \text{amod-adiv-equality}$
apply ($\text{case-tac } z \bmod a = 0, \text{ simp add:ant-0 aadd-0-r, blast,}$
 $\text{thin-tac } \{x. x \in v \text{ ' carrier } K \wedge 0 < x\} \subseteq \text{LBset } 1,$
 $\text{thin-tac } v w \neq \infty, \text{ thin-tac } v w \neq -\infty$)

apply ($\text{frule AMin-k[of } v], \text{erule bexE,}$
 drule sym,
 drule sym,
 drule sym,
 $\text{rotate-tac } -1, \text{drule sym})$

apply ($\text{cut-tac } z = z \text{ in } z\text{-in-aug-inf,}$
 $\text{cut-tac } z = (z \text{ div } a) \text{ and } x = a \text{ in } \text{asp-z-Z,}$
 $\text{cut-tac } z = z \bmod a \text{ in } z\text{-in-aug-inf,}$
 $\text{frule-tac } a = \text{ant } z \text{ and } b = (z \text{ div } a) *_a \text{ ant } a \text{ and}$
 $c = \text{ant } (z \bmod a) \text{ in } \text{ant-sol, assumption+},$
 $\text{subst asprod-mult, simp, assumption, simp,}$
 $\text{frule-tac } x = k \text{ and } z = z \text{ div } a \text{ in } \text{val-exp[of } v],$
 $(\text{erule conjE})+, \text{assumption, simp, simp,}$
 $\text{thin-tac } (z \text{ div } a) *_a v k = v (k_K^{(z \text{ div } a)}),$
 $\text{erule conjE})$

apply ($\text{frule-tac } x = k \text{ and } n = z \text{ div } a \text{ in } \text{field-potent-nonzero1,}$
 $\text{assumption+},$
 $\text{frule-tac } a = k \text{ and } n = z \text{ div } a \text{ in } \text{npowf-mem, assumption,}$
 $\text{frule-tac } x1 = k_K^{(z \text{ div } a)} \text{ in } \text{value-of-inv[THEN sym, of } v], \text{assumption+},$
 $\text{simp add:diff-ant-def,}$
 $\text{thin-tac } -v (k_K^{(z \text{ div } a)}) = v ((k_K^{(z \text{ div } a)})^{-K}),$
 $\text{cut-tac } x = k_K^{(z \text{ div } a)} \text{ in } \text{invf-closed1, simp,}$
 $\text{simp, erule conjE,}$
 $\text{frule-tac } x1 = w \text{ and } y1 = (k_K^{(z \text{ div } a)})^{-K} \text{ in}$
 $\text{val-t2p[THEN sym, of } v], \text{assumption+}, \text{simp,}$
 $\text{cut-tac field-is-ring,}$
 $\text{thin-tac } v w + v ((k_K^{(z \text{ div } a)})^{-K}) = \text{ant } (z \bmod a),$
 $\text{thin-tac } v (k_K^{(z \text{ div } a)}) + \text{ant } (z \bmod a) = v w,$
 $\text{frule-tac } x = w \text{ and } y = (k_K^{(z \text{ div } a)})^{-K} \text{ in}$
 $\text{Ring.ring-tOp-closed[of } K], \text{assumption+})$

apply ($\text{frule valuation-map[of } v],$
 $\text{frule-tac } a = w \cdot_r (k_K^{(z \text{ div } a)})^{-K} \text{ in } \text{mem-in-image[of } v$
 $\text{carrier } K Z_\infty], \text{assumption+}, \text{simp})$

apply ($\text{thin-tac AMin } \{x. x \in v \text{ ' carrier } K \wedge 0 < x\} = v k,$
 $\text{thin-tac } v \in \text{carrier } K \rightarrow Z_\infty,$
 $\text{subgoal-tac } 0 < v (w \cdot_r (k_K^{(z \text{ div } a)})^{-K}),$
 $\text{drule-tac } a = v (w \cdot_r (k_K^{(z \text{ div } a)})^{-K}) \text{ in } \text{forall-spec,}$
 $\text{simp add:image-def})$

apply ($\text{drule sym, simp})$
using $\text{not-zle pos-mod-bound apply blast}$
using $\text{pos-mod-sign zle-imp-zless-or-eq apply (metis Zero-ant-def aless)}$
done

lemma (*in Corps*) *val-principalTr1*: $\llbracket \text{valuation } K \ v \rrbracket \implies$
 $Lv \ K \ v \in v \text{ ' } (\text{carrier } K - \{0\}) \wedge$
 $(\forall w \in v \text{ ' } \text{carrier } K. \exists a. w = a * Lv \ K \ v) \wedge 0 < Lv \ K \ v$

apply (*rule conjI*,
frule val-Pg[*of v*], *erule conjE*,
simp add:image-def, *frule sym*, *thin-tac v* (*Pg K v*) = *Lv K v*,
erule conjE, *blast*)

apply (*rule conjI*,
rule ballI, *simp add:image-def*, *erule bexE*)

apply (
frule-tac x = x in value-in-aug-inf[*of v*], *assumption*,
frule sym, *thin-tac w = v x*, *simp add:aug-inf-def*,
cut-tac a = w in mem-ant, *simp*, *erule disjE*, *erule exE*,
frule-tac x = x in value-noninf-nonzero[*of v*], *assumption+*,
simp, *frule amin-generateTr*[*of v*])

apply (*drule-tac x = x in bspec*, *simp*,
erule exE,
frule AMin-z[*of v*], *erule exE*, *simp add:Lv-def*,
simp add:asprod-mult, *frule sym*, *thin-tac za * a = z*,
simp, *subst a-z-z*[*THEN sym*], *blast*)

apply (*simp add:Lv-def*,
frule AMin-z[*of v*], *erule exE*, *simp*,
frule Lv-pos[*of v*], *simp add:Lv-def*,
frule-tac m1 = a in a-i-pos[*THEN sym*], *blast*,
simp add:Lv-pos)

done

lemma (*in Corps*) *val-principalTr2*: $\llbracket \text{valuation } K \ v;$
 $c \in v \text{ ' } (\text{carrier } K - \{0\}) \wedge (\forall w \in v \text{ ' } \text{carrier } K. \exists a. w = a * c) \wedge 0 < c;$
 $d \in v \text{ ' } (\text{carrier } K - \{0\}) \wedge (\forall w \in v \text{ ' } \text{carrier } K. \exists a. w = a * d) \wedge 0 < d \rrbracket$
 $\implies c = d$

apply ((*erule conjE*)+,
drule-tac x = d in bspec,
simp add:image-def, *erule bexE*, *blast*,
drule-tac x = c in bspec,
simp add:image-def, *erule bexE*, *blast*)

apply ((*erule exE*)+,
drule sym, *simp*,
simp add:image-def, (*erule bexE*)+, *simp*,
(*erule conjE*)+,
frule-tac x = x in val-nonzero-z[*of v*], *assumption+*, *erule exE*,
frule-tac x = xa in val-nonzero-z[*of v*], *assumption+*, *erule exE*,
simp) **apply** (
subgoal-tac a $\neq \infty \wedge a \neq -\infty$, *subgoal-tac aa $\neq \infty \wedge aa \neq -\infty$* ,
cut-tac a = a in mem-ant, *cut-tac a = aa in mem-ant*, *simp*,

(erule exE)+, simp add:a-z-z,
 thin-tac c = ant z, frule sym, thin-tac zb * z = za, simp)
apply (subgoal-tac 0 < zb,
 cut-tac a = zc **and** b = zb **in** mult.commute, simp,
 simp add:pos-zmult-eq-1-iff,
 rule contrapos-pp, simp+,
 cut-tac x = 0 **and** y = zb **in** less-linear, simp,
 thin-tac $\neg 0 < zb$,
 erule disjE, simp,
 frule-tac i = 0 **and** j = z **and** k = zb **in** zmult-zless-mono-neg,
 assumption+, simp add:mult.commute)
apply (rule contrapos-pp, simp+, thin-tac $a \neq \infty \wedge a \neq -\infty$,
 erule disjE, simp, rotate-tac 5, drule sym,
 simp, simp, rotate-tac 5, drule sym, simp)
apply (rule contrapos-pp, simp+,
 erule disjE, simp, rotate-tac 4,
 drule sym, simp, simp,
 rotate-tac 4, drule sym,
 simp)
done

lemma (in Corps) val-principal:valuation K v $\implies$
 $\exists!x0. x0 \in v \text{ ' (carrier K - \{0\}) } \wedge$
 $(\forall w \in v \text{ ' (carrier K). } \exists (a::ant). w = a * x0) \wedge 0 < x0$
by (rule ex-ex1I,
 frule val-principalTr1[of v], blast,
 rule-tac c = x0 **and** d = y **in** val-principalTr2[of v],
 assumption+)

lemma (in Corps) n-val-defTr:[valuation K v; w $\in$ carrier K] $\implies$
 $\exists!a. a * Lv K v = v w$

apply (rule ex-ex1I,
 frule AMin-k[of v],
 frule Lv-pos[of v], simp add:Lv-def,
 erule bexE,
 frule-tac x = k **in** val-nonzero-z[of v], simp, simp,
 erule exE, simp, (erule conjE)+)
apply (case-tac w = 0_K, simp add:value-of-zero,
 frule-tac m = z **in** a-i-pos, blast)
apply (frule amin-generateTr[of v],
 drule-tac x = w **in** bspec, simp, simp)
apply (
 erule exE, simp add:asprod-mult,
 subst a-z-z[THEN sym], blast)
apply (frule AMin-k[of v]) **apply** (erule bexE,
 frule Lv-pos[of v], simp add:Lv-def) **apply** (
 erule conjE,
 frule-tac x = k **in** val-nonzero-z[of v], assumption+,
 erule exE, simp) **apply** (

case-tac $w = \mathbf{0}_K$, *simp* *del:a-i-pos* *add:value-of-zero*,
subgoal-tac $y = \infty$, *simp*, *rule* *contrapos-pp*, *simp+*,
cut-tac $a = a$ **in** *mem-ant*, *simp*,
erule *disjE*, *simp*, *erule* *exE*, *simp* *add:a-z-z*)
apply (*rule* *contrapos-pp*, *simp+*,
cut-tac $a = y$ **in** *mem-ant*, *simp*, *erule* *disjE*, *simp*,
erule *exE*, *simp* *add:a-z-z*,
frule-tac $x = w$ **in** *val-nonzero-z[of v]*, *assumption+*,
erule *exE*, *simp*, *cut-tac* $a = a$ **in** *mem-ant*,
erule *disjE*, *simp*, *frule* *sym*, *thin-tac* $-\infty = \text{ant } za$, *simp*,
erule *disjE*, *erule* *exE*, *simp* *add:a-z-z*)
apply (*cut-tac* $a = y$ **in** *mem-ant*,
erule *disjE*, *simp*, *rotate-tac* $\exists$, *drule* *sym*,
simp, *erule* *disjE*, *erule* *exE*, *simp* *add:a-z-z*, *frule* *sym*,
thin-tac $zb * z = za$, *simp*, *simp*,
rotate-tac $\exists$, *drule* *sym*,
simp, *simp*, *frule* *sym*, *thin-tac* $\infty = \text{ant } za$, *simp*)
done

lemma (**in** *Corps*) *n-valTr*: $\llbracket \text{valuation } K v; x \in \text{carrier } K \rrbracket \implies$
 $(\text{THE } l. (l * (Lv K v)) = v x) * (Lv K v) = v x$
by (*rule* *theI'*, *rule* *n-val-defTr*, *assumption+*)

lemma (**in** *Corps*) *n-val*: $\llbracket \text{valuation } K v; x \in \text{carrier } K \rrbracket \implies$
 $(n\text{-val } K v x) * (Lv K v) = v x$
by (*frule* *n-valTr[of v x]*, *assumption+*, *simp* *add:n-val-def*)

lemma (**in** *Corps*) *val-pos-n-val-pos*: $\llbracket \text{valuation } K v; x \in \text{carrier } K \rrbracket \implies$
 $(0 \leq v x) = (0 \leq n\text{-val } K v x)$
apply (*frule* *n-val[of v x]*, *assumption+*,
drule *sym*,
frule *Lv-pos[of v]*,
frule *Lv-z[of v]*, *erule* *exE*, *simp*)
apply (*frule-tac* $w = z$ **and** $x = 0$ **and** $y = n\text{-val } K v x$ **in** *amult-pos-mono-r*,
simp *add:amult-0-l*)
done

lemma (**in** *Corps*) *n-val-in-aug-inf*: $\llbracket \text{valuation } K v; x \in \text{carrier } K \rrbracket \implies$
 $n\text{-val } K v x \in Z_\infty$
apply (*cut-tac* *field-is-ring*, *frule* *Ring.ring-zero[of K]*,
frule *Lv-pos[of v]*,
frule *Lv-z[of v]*, *erule* *exE*,
simp *add:aug-inf-def*)
apply (*rule* *contrapos-pp*, *simp+*)
apply (*case-tac* $x = \mathbf{0}_K$, *simp*,
frule *n-val[of v 0]*,
simp *add:value-of-zero*, *simp* *add:value-of-zero*)
apply (*frule* *n-val[of v x]*, *simp*,

$\text{frule val-nonzero-z[of } v \ x], \text{ assumption+},$
 $\text{erule exE, simp, rotate-tac } -2, \text{ drule sym,}$
 $\text{simp})$
done

lemma (in Corps) $n\text{-val-0}:\llbracket \text{valuation } K \ v; x \in \text{carrier } K; v \ x = 0 \rrbracket$
 $\implies n\text{-val } K \ v \ x = 0$
by (frule Lv-z[of v], erule exE,
 $\text{frule Lv-pos[of } v],$
 $\text{frule n-val[of } v \ x], \text{ simp, simp,}$
 $\text{rule-tac } z = z \text{ and } a = n\text{-val } K \ v \ x \text{ in } a\text{-a-z-0, assumption+})$

lemma (in Corps) $\text{value-n0-n-val-n0}:\llbracket \text{valuation } K \ v; x \in \text{carrier } K; v \ x \neq 0 \rrbracket \implies$
 $n\text{-val } K \ v \ x \neq 0$
apply (frule n-val[of v x],
 $\text{rule contrapos-pp, simp+}, \text{frule Lv-z[of } v],$
 $\text{erule exE, simp, simp only:ant-0[THEN sym]})$
apply (rule contrapos-pp, simp+,
 $\text{simp add:a-z-z})$
done

lemma (in Corps) $\text{val-0-n-val-0}:\llbracket \text{valuation } K \ v; x \in \text{carrier } K \rrbracket \implies$
 $(v \ x = 0) = (n\text{-val } K \ v \ x = 0)$
apply (rule iffI,
 $\text{simp add:n-val-0})$
apply (rule contrapos-pp, simp+,
 $\text{frule value-n0-n-val-n0[of } v \ x], \text{ assumption+})$
apply simp
done

lemma (in Corps) $\text{val-noninf-n-val-noninf}:\llbracket \text{valuation } K \ v; x \in \text{carrier } K \rrbracket \implies$
 $(v \ x \neq \infty) = (n\text{-val } K \ v \ x \neq \infty)$
by (frule Lv-z[of v], erule exE,
 $\text{frule Lv-pos[of } v], \text{ simp,}$
 $\text{frule n-val[THEN sym, of } v \ x], \text{simp, simp,}$
 $\text{thin-tac } v \ x = n\text{-val } K \ v \ x * \text{ant } z,$
 $\text{rule iffI, rule contrapos-pp, simp+},$
 $\text{cut-tac mem-ant[of } n\text{-val } K \ v \ x], \text{erule disjE, simp,}$
 $\text{erule disjE, erule exE, simp add:a-z-z, simp, simp})$

lemma (in Corps) $\text{val-inf-n-val-inf}:\llbracket \text{valuation } K \ v; x \in \text{carrier } K \rrbracket \implies$
 $(v \ x = \infty) = (n\text{-val } K \ v \ x = \infty)$
by (cut-tac val-noninf-n-val-noninf[of v x], simp, assumption+)

lemma (in Corps) $\text{val-eq-n-val-eq}:\llbracket \text{valuation } K \ v; x \in \text{carrier } K; y \in \text{carrier } K \rrbracket$
 $\implies (v \ x = v \ y) = (n\text{-val } K \ v \ x = n\text{-val } K \ v \ y)$
apply (subst n-val[THEN sym, of v x], assumption+,
 $\text{subst n-val[THEN sym, of } v \ y], \text{ assumption+},$
 $\text{frule Lv-pos[of } v], \text{frule Lv-z[of } v], \text{erule exE, simp,}$

$\text{frule-tac } s = z \text{ in } \text{zless-neq}[\text{THEN not-sym, of } 0])$
apply (rule iffI)
apply (rule-tac $z = z$ in amult-eq-eq-r[of - n-val $K \ v \ x$ n-val $K \ v \ y$],
 assumption+)
apply simp
done

lemma (in Corps) val-poss-n-val-poss:[[valuation $K \ v$; $x \in \text{carrier } K$]] $\implies$
 $(0 < v \ x) = (0 < \text{n-val } K \ v \ x)$
apply (simp add:less-le,
 frule val-pos-n-val-pos[of $v \ x$], assumption+,
 rule iffI, erule conjE, simp,
 simp add:value-n0-n-val-n0[of $v \ x$])
apply (drule sym,
 erule conjE, simp,
 frule-tac val-0-n-val-0[THEN sym, of $v \ x$], assumption+,
 simp)
done

lemma (in Corps) n-val-Pg:valuation $K \ v \implies \text{n-val } K \ v \ (\text{Pg } K \ v) = 1$
apply (frule val-Pg[of v], simp, (erule conjE)+,
 frule n-val[of $v \ \text{Pg } K \ v$], simp, frule Lv-z[of v], erule exE, simp,
 frule Lv-pos[of v], simp, frule-tac $i = 0$ and $j = z$ in zless-neq)
apply (rotate-tac -1, frule not-sym, thin-tac $0 \neq z$,
 subgoal-tac $\text{n-val } K \ v \ (\text{Pg } K \ v) * \text{ant } z = 1 * \text{ant } z$,
 rule-tac $z = z$ in adiv-eq[of - n-val $K \ v \ (\text{Pg } K \ v) \ 1$], assumption+,
 simp add:amult-one-l)
done

lemma (in Corps) n-val-valuationTr1:valuation $K \ v \implies$
 $\forall x \in \text{carrier } K. \text{n-val } K \ v \ x \in Z_\infty$
by (rule ballI,
 frule n-val[of v], assumption,
 frule-tac $x = x$ in value-in-aug-inf[of v], assumption,
 frule Lv-pos[of v], simp add:aug-inf-def,
 frule Lv-z[of v], erule exE, simp,
 rule contrapos-pp, simp+)

lemma (in Corps) n-val-t2p:[[valuation $K \ v$; $x \in \text{carrier } K$; $y \in \text{carrier } K$]] $\implies$
 $\text{n-val } K \ v \ (x \cdot_r y) = \text{n-val } K \ v \ x + (\text{n-val } K \ v \ y)$
apply (cut-tac field-is-ring,
 frule Ring.ring-tOp-closed[of $K \ x \ y$], assumption+,
 frule n-val[of $v \ x \cdot_r y$], assumption+,
 frule Lv-pos[of v],
 simp add:val-t2p,
 frule n-val[THEN sym, of $v \ x$], assumption+,
 frule n-val[THEN sym, of $v \ y$], assumption+, simp,
 frule Lv-z[of v], erule exE, simp)
apply (subgoal-tac ant $z \neq 0$)

apply (*frule-tac* $z1 = z$ **in** *amult-distrib1* [*THEN sym*, *of - n-val K v x*
n-val K v y], *simp*,
thin-tac *n-val K v x* * *ant z* + *n-val K v y* * *ant z* =
 (*n-val K v x* + *n-val K v y*) * *ant z*,
rule-tac $z = z$ **and** $a = n\text{-val } K \ v \ (x \cdot_r y)$ **and**
 $b = n\text{-val } K \ v \ x + n\text{-val } K \ v \ y$ **in** *adiv-eq*, *simp*, *assumption+*, *simp*)
done

lemma (**in** *Corps*) *n-val-valuationTr2*: [*valuation K v*; $x \in \text{carrier } K$;
 $y \in \text{carrier } K$] $\implies$
 $\text{amin } (n\text{-val } K \ v \ x) \ (n\text{-val } K \ v \ y) \leq (n\text{-val } K \ v \ (x \pm y))$
apply (*frule* *n-val* [*THEN sym*, *of v x*], *assumption+*,
frule *n-val* [*THEN sym*, *of v y*], *assumption+*,
frule *n-val* [*THEN sym*, *of v x ± y*],
cut-tac *field-is-ring*, *frule* *Ring.ring-is-ag*[*of K*],
rule *aGroup.ag-pOp-closed*, *assumption+*)
apply (*frule* *amin-le-plus*[*of v x y*], *assumption+*, *simp*,
simp *add:amult-commute*[*of - Lv K v*],
frule *Lv-z*[*of v*], *erule* *exE*, *simp*,
frule *Lv-pos*[*of v*], *simp*,
simp *add:amin-amult-pos*, *simp* *add:amult-pos-mono-l*)
done

lemma (**in** *Corps*) *n-val-valuation:valuation K v* $\implies$
 $\text{valuation } K \ (n\text{-val } K \ v)$
apply (*cut-tac* *field-is-ring*, *frule* *Ring.ring-is-ag*)
apply (*frule* *Lv-z*[*of v*], *erule* *exE*, *frule* *Lv-pos*[*of v*], *simp*,
subst *valuation-def*)
apply (*rule* *conjI*, *simp* *add:n-val-def restrict-def extensional-def*)
apply (*rule* *conjI*, *simp* *add:n-val-valuationTr1*)
apply (*rule* *conjI*, *frule* *n-val*[*of v 0*],
simp *add:Ring.ring-zero*,
frule *Lv-z*[*of v*], *erule* *exE*, *frule* *Lv-pos*[*of v*],
cut-tac *mem-ant*[*of n-val K v (0)*], *erule* *disjE*,
simp *add:value-of-zero*,
erule *disjE*, *erule* *exE*, *simp* *add:a-z-z value-of-zero*, *assumption+*)
apply (*rule* *conjI*, *rule* *ballI*,
frule-tac $x = x$ **in** *val-nonzero-noninf*[*of v*], *simp+*,
simp *add:val-noninf-n-val-noninf*)
apply (*rule* *conjI*, (*rule* *ballI*) $+$, *simp* *add:n-val-t2p*,
rule *conjI*, *rule* *ballI*, *rule* *impI*,
frule *Lv-z*[*of v*], *erule* *exE*,
frule *Lv-pos*[*of v*], *simp*,
frule-tac $x = x$ **in** *n-val*[*of v*], *simp*,
frule-tac $w1 = z$ **and** $x1 = 0$ **and** $y1 = n\text{-val } K \ v \ x$ **in**
amult-pos-mono-r [*THEN sym*], *simp* *add:amult-0-l*,
frule-tac $x = x$ **in** *val-axiom4*[*of v*], *assumption+*,
frule-tac $x1 = 1_r \pm x$ **in** *n-val* [*THEN sym*, *of v*],
frule *Ring.ring-is-ag*[*of K*],

$\text{rule } aGroup.ag-pOp-closed, \text{ assumption+}, \text{ simp add: Ring.ring-one},$
 $\text{assumption},$
 $\text{frule-tac } w = z \text{ and } x = 0 \text{ and } y = n\text{-val } K \ v \ (1_r \pm x)$
 $\text{in } amult\text{-pos-mono-r},$
 $\text{simp add: amult-0-l})$

apply ($\text{frule val-axiom5[of } v], \text{erule exE},$
 $(\text{erule conjE})+,$
 $\text{frule-tac } x = x \text{ in } value\text{-n0-n-val-n0[of } v], \text{assumption+},$
 $\text{frule-tac } x = x \text{ in } val\text{-noninf-n-val-noninf}, \text{simp},$
 $\text{blast})$

done

lemma (**in** $Corps$) $n\text{-val-le-val}:\llbracket valuation \ K \ v; x \in carrier \ K; 0 \leq (v \ x) \rrbracket \implies$
 $(n\text{-val } K \ v \ x) \leq (v \ x)$

by ($\text{subst } n\text{-val}[THEN \ sym, \text{ of } v \ x], \text{assumption+},$
 $\text{frule } Lv\text{-pos[of } v],$
 $\text{simp add: val-pos-n-val-pos[of } v \ x],$
 $\text{frule } Lv\text{-z[of } v], \text{erule exE},$
 $\text{cut-tac } b = z \text{ and } x = n\text{-val } K \ v \ x \text{ in } amult\text{-pos}, \text{simp+},$
 $\text{simp add: asprod-amult}, \text{simp add: amult-commute})$

lemma (**in** $Corps$) $n\text{-val-surj:valuation } K \ v \implies$
 $\exists x \in carrier \ K. n\text{-val } K \ v \ x = 1$

apply ($\text{frule } Lv\text{-z[of } v], \text{erule exE},$
 $\text{frule } Lv\text{-pos[of } v],$
 $\text{frule } AMin\text{-k[of } v], \text{erule bexE}, \text{frule-tac } x = k \text{ in } n\text{-val[of } v], \text{simp},$
 $\text{simp add: Lv-def})$

apply ($\text{subgoal-tac } n\text{-val } K \ v \ k * ant \ z = 1 * ant \ z,$
 $\text{subgoal-tac } z \neq 0,$
 $\text{frule-tac } z = z \text{ and } a = n\text{-val } K \ v \ k \text{ and } b = 1 \text{ in } amult\text{-eq-eq-r},$
 $\text{assumption}, \text{blast}, \text{simp}, \text{simp add: amult-one-l})$

done

lemma (**in** $Corps$) $n\text{-value-in-aug-inf}:\llbracket valuation \ K \ v; x \in carrier \ K \rrbracket \implies$
 $n\text{-val } K \ v \ x \in Z_\infty$

by ($\text{frule } n\text{-val[of } v \ x], \text{assumption},$
 $\text{simp add: aug-inf-def}, \text{rule contrapos-pp}, \text{simp+},$
 $\text{frule } Lv\text{-pos[of } v], \text{frule } Lv\text{-z[of } v], \text{erule exE}, \text{simp},$
 $\text{frule value-in-aug-inf[of } v \ x], \text{assumption+}, \text{simp add: aug-inf-def})$

lemma (**in** $Corps$) $val\text{-surj-n-valTr}:\llbracket valuation \ K \ v; \exists x \in carrier \ K. v \ x = 1 \rrbracket$
 $\implies Lv \ K \ v = 1$

apply ($\text{erule bexE},$
 $\text{frule-tac } x = x \text{ in } n\text{-val[of } v],$
 $\text{simp}, \text{frule } Lv\text{-pos[of } v])$

apply ($\text{frule-tac } w = Lv \ K \ v \text{ and } x = n\text{-val } K \ v \ x \text{ in } amult\text{-1-both})$

apply simp+

done

```

lemma (in Corps) val-surj-n-val:[valuation K v;  $\exists x \in \text{carrier } K. v \ x = 1$ ]  $\implies$ 
    (n-val K v) = v
apply (rule funcset-eq[of - carrier K],
    simp add:n-val-def restrict-def extensional-def,
    simp add:valuation-def)
apply (rule ballI,
    frule val-surj-n-valTr[of v], assumption+,
    frule-tac x = x in n-val[of v], assumption+,
    simp add:amult-one-r)
done

lemma (in Corps) n-val-n-val:valuation K v  $\implies$ 
    n-val K (n-val K v) = n-val K v
by (frule n-val-valuation[of v],
    frule n-val-surj[of v],
    simp add:val-surj-n-val)

lemma nnonzero-annnonzero:0 < N  $\implies$  an N  $\neq$  0
apply (simp only:an-0[THEN sym])
apply (subst aneq-natneq, simp)
done

```

2.3 Valuation ring

definition

$Vr :: [(r, 'm) \text{ Ring-scheme}, 'r \Rightarrow \text{ant}] \Rightarrow (r, 'm) \text{ Ring-scheme}$ **where**
 $Vr \ K \ v = Sr \ K \ (\{x. x \in \text{carrier } K \wedge 0 \leq (v \ x)\})$

definition

$vp :: [(r, 'm) \text{ Ring-scheme}, 'r \Rightarrow \text{ant}] \Rightarrow 'r \text{ set}$ **where**
 $vp \ K \ v = \{x. x \in \text{carrier } (Vr \ K \ v) \wedge 0 < (v \ x)\}$

definition

$r\text{-apow} :: [(r, 'm) \text{ Ring-scheme}, 'r \text{ set}, \text{ant}] \Rightarrow 'r \text{ set}$ **where**
 $r\text{-apow } R \ I \ a = (\text{if } a = \infty \text{ then } \{0_R\} \text{ else}$
 $(\text{if } a = 0 \text{ then } \text{carrier } R \text{ else } I \diamond R \ (na \ a)))$

abbreviation

$RAPOW \ (\langle 3- \ - \ \rangle) [62,62,63]62$ **where**
 $I^R \ a == r\text{-apow } R \ I \ a$

lemma (**in** Ring) ring-pow-apow:ideal R I $\implies$
 $I \diamond R \ n = I^R \ (an \ n)$

apply (simp add:r-apow-def)
apply (case-tac n = 0, simp)
apply (simp add:nnonzero-annnonzero)

apply (*simp add:an-neq-inf na-an*)
done

lemma (**in** *Ring*) *r-apow-Suc:ideal R I $\implies I^R (an (Suc 0)) = I$*
apply (*cut-tac an-1, simp add:r-apow-def*)
apply (*simp add:a0-neq-1[THEN not-sym]*)
 apply (*simp only:ant-1[THEN sym]*)
 apply (*simp del:ant-1 add:z-neq-inf[of 1, THEN not-sym]*)
 apply (*simp add:na-1*)
 apply (*simp add:idealprod-whole-r*)
done

lemma (**in** *Ring*) *apow-ring-pow:ideal R I $\implies$*
 $I \diamond R \ n = I^R (an \ n)$
apply (*simp add:r-apow-def*)
apply (*case-tac n = 0, simp add:an-0*)
apply (*simp add: aless-nat-less[THEN sym],*
 cut-tac an-neq-inf[of n],
 simp add: less-le[of 0 an n] na-an)
done

lemma (**in** *Corps*) *Vr-ring:valuation K v $\implies$ Ring (Vr K v)*
apply (*cut-tac field-is-ring, frule Ring.ring-is-ag[of K],*
 simp add: Vr-def, rule Ring.Sr-ring, assumption+)
 apply (*simp add:sr-def*)
 apply (*intro conjI subsetI*)
apply (*simp-all add: value-of-one Ring.ring-one[of K]*)
apply ((*rule allI, rule impI*)+,
 (*erule conjE*)+, *rule conjI, rule aGroup.ag-pOp-closed, assumption+*,
 rule aGroup.ag-mOp-closed, assumption+)
apply (*frule-tac x = x and y = -_a y in amin-le-plus[of v], assumption+*,
 rule aGroup.ag-mOp-closed, assumption+,
 simp add:val-minus-eq[of v]) **apply** (
 frule-tac z = 0 and x = v x and y = v y in amin-ge1, assumption+,
 frule-tac i = 0 and j = amin (v x) (v y) and k = v (x $\pm$ -_a y) in
 ale-trans, assumption+, simp)
by (*simp add: Ring.ring-tOp-closed aadd-two-pos val-t2p*)

lemma (**in** *Corps*) *val-pos-mem-Vr: [valuation K v; x $\in$ carrier K] $\implies$*
 $(0 \leq (v \ x)) = (x \in \text{carrier } (Vr \ K \ v))$
by (*rule iffI, (simp add: Vr-def Sr-def)+*)

lemma (**in** *Corps*) *val-poss-mem-Vr: [valuation K v; x $\in$ carrier K; 0 < (v x)]*
 $\implies x \in \text{carrier } (Vr \ K \ v)$
by (*frule aless-imp-le[of 0 v x], simp add: val-pos-mem-Vr*)

lemma (**in** *Corps*) *Vr-one:valuation K v $\implies 1_{rK} \in \text{carrier } (Vr \ K \ v)$*
by (*cut-tac field-is-ring, frule Ring.ring-one[of K],*
 frule val-pos-mem-Vr[of v 1_r], assumption+,

simp add:value-of-one)

lemma (in Corps) Vr-mem-f-mem: $\llbracket \text{valuation } K \ v; x \in \text{carrier } (Vr \ K \ v) \rrbracket$
 $\implies x \in \text{carrier } K$
by (*simp add: Vr-def Sr-def*)

lemma (in Corps) Vr-0-f-0: $\text{valuation } K \ v \implies \mathbf{0}_{Vr \ K \ v} = \mathbf{0}$
by (*simp add: Vr-def Sr-def*)

lemma (in Corps) Vr-1-f-1: $\text{valuation } K \ v \implies 1_r(Vr \ K \ v) = 1_r$
by (*simp add: Vr-def Sr-def*)

lemma (in Corps) Vr-pOp-f-pOp: $\llbracket \text{valuation } K \ v; x \in \text{carrier } (Vr \ K \ v);$
 $y \in \text{carrier } (Vr \ K \ v) \rrbracket \implies x \pm_{Vr \ K \ v} y = x \pm y$
by (*simp add: Vr-def Sr-def*)

lemma (in Corps) Vr-mOp-f-mOp: $\llbracket \text{valuation } K \ v; x \in \text{carrier } (Vr \ K \ v) \rrbracket$
 $\implies -_a(Vr \ K \ v) \ x = -_a \ x$
by (*simp add: Vr-def Sr-def*)

lemma (in Corps) Vr-tOp-f-tOp: $\llbracket \text{valuation } K \ v; x \in \text{carrier } (Vr \ K \ v);$
 $y \in \text{carrier } (Vr \ K \ v) \rrbracket \implies x \cdot_r(Vr \ K \ v) \ y = x \cdot_r y$
by (*simp add: Vr-def Sr-def*)

lemma (in Corps) Vr-pOp-le: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K;$
 $y \in \text{carrier } (Vr \ K \ v) \rrbracket \implies v \ x \leq (v \ x + (v \ y))$
apply (*frule val-pos-mem-Vr[THEN sym, of v y],*
simp add: Vr-mem-f-mem, simp, frule aadd-pos-le[of v y v x],
simp add: aadd-commute)
done

lemma (in Corps) Vr-integral: $\text{valuation } K \ v \implies \text{Idomain } (Vr \ K \ v)$
apply (*simp add: Idomain-def,*
simp add: Vr-ring, simp add: Idomain-axioms-def,
rule allI, rule impI, rule allI, (rule impI)+,
simp add: Vr-tOp-f-tOp, simp add: Vr-0-f-0)
apply (*rule contrapos-pp, simp+, erule conjE,*
cut-tac field-is-idom,
frule-tac x = a in Vr-mem-f-mem[of v], assumption,
frule-tac x = b in Vr-mem-f-mem[of v], assumption,
frule-tac x = a and y = b in Idomain.idom-tOp-nonzeros[of K],
assumption+, simp)
done

lemma (in Corps) Vr-exp-mem: $\llbracket \text{valuation } K \ v; x \in \text{carrier } (Vr \ K \ v) \rrbracket$
 $\implies x^{\sim K \ n} \in \text{carrier } (Vr \ K \ v)$
by (*frule Vr-ring[of v],*
induct-tac n, simp add: Vr-one,
simp add: Vr-tOp-f-tOp[THEN sym, of v],

simp add: Ring.ring-tOp-closed)

lemma (in Corps) *Vr-exp-f-exp*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } (Vr \ K \ v) \rrbracket \implies$

$$x \frown (Vr \ K \ v) \ n = x \frown^K n$$

apply (*induct-tac* *n*,
simp, *simp add: Vr-1-f-1*, *simp*,
thin-tac $x \frown (Vr \ K \ v) \ n = x \frown^K n$)

apply (*rule Vr-tOp-f-tOp*, *assumption+*,
simp add: Vr-exp-mem, *assumption*)

done

lemma (in Corps) *Vr-potent-nonzero*: $\llbracket \text{valuation } K \ v;$
 $x \in \text{carrier } (Vr \ K \ v) - \{0_{Vr \ K \ v}\} \rrbracket \implies x \frown^K n \neq 0_{Vr \ K \ v}$

apply (*frule Vr-mem-f-mem*[*of v x*], *simp*,
simp add: Vr-0-f-0, *erule conjE*)

apply (*frule Vr-mem-f-mem*[*of v x*], *assumption+*,
simp add: field-potent-nonzero)

done

lemma (in Corps) *elem-0-val-if*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; v \ x = 0 \rrbracket$
 $\implies x \in \text{carrier } (Vr \ K \ v) \wedge x \frown^K \in \text{carrier } (Vr \ K \ v)$

apply (*frule val-pos-mem-Vr*[*of v x*], *assumption*, *simp*)

apply (*frule value-zero-nonzero*[*of v x*], *simp add: Vr-mem-f-mem*, *simp*)

apply (*frule value-of-inv*[*of v x*], *assumption+*,
simp, *subst val-pos-mem-Vr*[*THEN sym*, *of v x* $\frown^K$], *assumption+*,
cut-tac invf-closed[*of x*], *simp+*)

done

lemma (in Corps) *elem0val*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; x \neq 0 \rrbracket \implies$
 $(v \ x = 0) = (x \in \text{carrier } (Vr \ K \ v) \wedge x \frown^K \in \text{carrier } (Vr \ K \ v))$

apply (*rule iffI*, *rule elem-0-val-if*[*of v*], *assumption+*,
erule conjE)

apply (*simp add: val-pos-mem-Vr*[*THEN sym*, *of v x*],
frule Vr-mem-f-mem[*of v x* $\frown^K$], *assumption+*,
simp add: val-pos-mem-Vr[*THEN sym*, *of v x* $\frown^K$],
simp add: value-of-inv, *frule ale-minus*[*of 0 - v x*],
simp add: a-minus-minus)

done

lemma (in Corps) *ideal-inc-elem0val-whole*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K;$
 $v \ x = 0; \text{ideal } (Vr \ K \ v) \ I; x \in I \rrbracket \implies I = \text{carrier } (Vr \ K \ v)$

apply (*frule elem-0-val-if*[*of v x*], *assumption+*, *erule conjE*,
frule value-zero-nonzero[*of v x*], *assumption+*,
frule Vr-ring[*of v*],
frule-tac $I = I$ **and** $x = x$ **and** $r = x \frown^K$ **in**
Ring.ideal-ring-multiple[*of Vr K v*], *assumption+*,
cut-tac invf-closed1[*of x*], *simp+*, (*erule conjE*) $+$)

apply (*simp add: Vr-tOp-f-tOp*, *cut-tac invf-inv*[*of x*], *simp+*,

$\text{simp add: Vr-1-f-1[THEN sym, of v],}$
 $\text{simp add: Ring.ideal-inc-one, simp+)}$
done

lemma (in Corps) $\text{vp-mem-Vr-mem:} \llbracket \text{valuation } K \ v; x \in (\text{vp } K \ v) \rrbracket \implies$
 $x \in \text{carrier } (Vr \ K \ v)$
by (rule val-poss-mem-Vr[of v x], assumption+, (simp add: vp-def
Vr-def Sr-def)+)

lemma (in Corps) $\text{vp-mem-val-poss:} \llbracket \text{valuation } K \ v; x \in \text{carrier } K \rrbracket \implies$
 $(x \in \text{vp } K \ v) = (0 < (v \ x))$
by (simp add: vp-def, simp add: Vr-def Sr-def less-ant-def)

lemma (in Corps) $\text{Pg-in-Vr: valuation } K \ v \implies \text{Pg } K \ v \in \text{carrier } (Vr \ K \ v)$
by (frule val-Pg[of v], erule conjE,
frule Lv-pos[of v], drule sym,
simp, erule conjE,
simp add: val-poss-mem-Vr)

lemma (in Corps) $\text{vp-ideal: valuation } K \ v \implies \text{ideal } (Vr \ K \ v) \ (\text{vp } K \ v)$
apply (cut-tac field-is-ring,
frule Vr-ring[of v],
rule Ring.ideal-condition1, assumption+,
rule subsetI, simp add: vp-mem-Vr-mem,
simp add: vp-def)
apply (frule val-Pg[of v],
frule Lv-pos[of v], simp, (erule conjE)+,
drule sym, simp,
frule val-poss-mem-Vr[of v Pg K v], assumption+, blast)

apply ((rule ballI)+,
frule-tac $x = x$ in vp-mem-Vr-mem[of v], assumption) **apply** (
frule-tac $x = y$ in vp-mem-Vr-mem[of v], assumption,
simp add: vp-def,
frule Ring.ring-is-ag[of Vr K v],
frule-tac $x = x$ and $y = y$ in aGroup.ag-pOp-closed, assumption+, simp)
apply (simp add: Vr-pOp-f-pOp,
cut-tac $x = v \ x$ and $y = v \ y$ in amin-le-l,
frule-tac $x = x$ and $y = y$ in amin-le-plus,
(simp add: Vr-mem-f-mem)+,
(frule-tac $z = 0$ and $x = v \ x$ and $y = v \ y$ in amin-gt, assumption+),
rule-tac $x = 0$ and $y = \text{amin } (v \ x) \ (v \ y)$ and $z = v \ (x \pm y)$ in
less-le-trans, assumption+)

apply ((rule ballI)+,
frule-tac $x1 = r$ in val-pos-mem-Vr[THEN sym, of v],
simp add: Vr-mem-f-mem, simp,
frule-tac $x = x$ in vp-mem-Vr-mem[of v], simp add: Vr-pOp-f-pOp,
simp add: vp-def, simp add: Ring.ring-tOp-closed,
simp add: Vr-tOp-f-tOp)

apply (*frule-tac* $x = r$ **in** *Vr-mem-f-mem*[*of v*], *assumption*+,
frule-tac $x = x$ **in** *Vr-mem-f-mem*[*of v*], *assumption*+,
simp add:val-t2p, *simp add:aadd-pos-poss*)
done

lemma (**in** *Corps*) *vp-not-whole:valuation* $K v \implies$
 $(vp\ K\ v) \neq carrier\ (Vr\ K\ v)$
apply (*cut-tac field-is-ring*, *frule Ring.ring-is-ag*[*of K*],
frule Vr-ring[*of v*])
apply (*rule contrapos-pp*, *simp*+,
drule sym,
frule Ring.ring-one[*of Vr K v*], *simp*,
simp add:Vr-1-f-1,
frule Ring.ring-one[*of K*])
apply (*simp only:vp-mem-val-poss*[*of v 1_r*],
simp add:value-of-one)
done

lemma (**in** *Ring*) *elem-out-ideal-nonzero*: $\llbracket ideal\ R\ I; x \in carrier\ R;$
 $x \notin I \rrbracket \implies x \neq \mathbf{0}_R$
by (*rule contrapos-pp*, *simp*+, *frule ideal-zero*[*of I*],
simp)

lemma (**in** *Corps*) *vp-prime:valuation* $K v \implies prime-ideal\ (Vr\ K\ v)\ (vp\ K\ v)$
apply (*simp add:prime-ideal-def*, *simp add:vp-ideal*)
apply (*rule conjI*)

apply (*rule contrapos-pp*, *simp*+,
frule Vr-ring[*of v*],
frule vp-ideal[*of v*],
frule Ring.ideal-inc-one[*of Vr K v vp K v*], *assumption*+,
simp add:vp-not-whole[*of v*])

apply ((*rule ballI*)+, *rule impI*, *rule contrapos-pp*, *simp*+, (*erule conjE*)+,
frule Vr-ring[*of v*]) **apply** (
frule-tac $x = x$ **in** *Vr-mem-f-mem*[*of v*], *assumption*) **apply** (
frule-tac $x = y$ **in** *Vr-mem-f-mem*[*of v*], *assumption*) **apply** (
frule vp-ideal[*of v*],
frule-tac $x = x$ **in** *Ring.elem-out-ideal-nonzero*[*of Vr K v vp K v*],
assumption+) **apply** (
frule-tac $x = y$ **in** *Ring.elem-out-ideal-nonzero*[*of Vr K v vp K v*],
assumption+, *simp add:Vr-0-f-0*,
simp add:Vr-tOp-f-tOp) **apply** (
frule-tac $x = x \cdot_r y$ **in** *vp-mem-val-poss*[*of v*],
cut-tac field-is-ring, *simp add:Ring.ring-tOp-closed*, *simp*)
apply (*cut-tac field-is-ring*,
frule-tac $x = x$ **and** $y = y$ **in** *Ring.ring-tOp-closed*, *assumption*+,
simp add:Ring.ring-tOp-closed[*of Vr K v*],

$\text{simp add:vp-def, simp add:aneg-less,}$
 $\text{frule-tac } x1 = x \text{ in val-pos-mem-Vr[THEN sym, of v], assumption+,}$
 $\text{frule-tac } x1 = y \text{ in val-pos-mem-Vr[THEN sym, of v], assumption+,}$
 $\text{frule-tac } P = x \in \text{carrier } (Vr K v) \text{ and } Q = 0 \leq v x \text{ in eq-prop,}$
 assumption,
 $\text{frule-tac } P = y \in \text{carrier } (Vr K v) \text{ and } Q = 0 \leq v y \text{ in eq-prop,}$
 assumption,
 $\text{frule-tac } x = v x \text{ and } y = 0 \text{ in ale-antisym, assumption+,}$
 $\text{frule-tac } x = v y \text{ and } y = 0 \text{ in ale-antisym, assumption+,}$
 $\text{simp add:val-t2p aadd-0-l)}$
done

lemma (in Corps) $\text{vp-pow-ideal:valuation } K v \implies$
 $\text{ideal } (Vr K v) ((vp K v) \diamond (Vr K v) n)$
by (frule Vr-ring[of v], frule vp-ideal[of v],
 $\text{simp add:Ring.ideal-pow-ideal})$

lemma (in Corps) $\text{vp-apow-ideal:} \llbracket \text{valuation } K v; 0 \leq n \rrbracket \implies$
 $\text{ideal } (Vr K v) ((vp K v) (Vr K v) n)$
apply (frule Vr-ring[of v])
apply (case-tac $n = 0$,
 $\text{simp add:r-apow-def, simp add:Ring.whole-ideal[of Vr K v]})$
apply (case-tac $n = \infty$,
 $\text{simp add:r-apow-def, simp add:Ring.zero-ideal})$
apply (simp add:r-apow-def, simp add:vp-pow-ideal)
done

lemma (in Corps) $\text{mem-vp-apow-mem-Vr:} \llbracket \text{valuation } K v;$
 $0 \leq N; x \in vp K v (Vr K v) N \rrbracket \implies x \in \text{carrier } (Vr K v)$
by (frule Vr-ring[of v], frule vp-apow-ideal[of v N], assumption,
 $\text{simp add:Ring.ideal-subset})$

lemma (in Corps) $\text{elem-out-vp-unit:} \llbracket \text{valuation } K v; x \in \text{carrier } (Vr K v);$
 $x \notin vp K v \rrbracket \implies v x = 0$
by (metis Vr-mem-f-mem ale-antisym aneg-le val-pos-mem-Vr vp-mem-val-poss)

lemma (in Corps) $\text{vp-maximal:valuation } K v \implies$
 $\text{maximal-ideal } (Vr K v) (vp K v)$
apply (frule Vr-ring[of v],
 $\text{simp add:maximal-ideal-def, simp add:vp-ideal})$

apply (frule vp-not-whole[of v],
 $\text{rule conjI, rule contrapos-pp, simp+, frule vp-ideal[of v],}$
 $\text{frule Ring.ideal-inc-one[of Vr K v vp K v], assumption+})$
apply simp

apply (rule equalityI,
 $\text{rule subsetI, simp, erule conjE,}$
 $\text{case-tac } x = vp K v, \text{simp, simp, rename-tac } X)$

```

apply (frule-tac  $A = X$  in sets-not-eq[of -  $vp\ K\ v$ ], assumption+,
  erule bexE,
  frule-tac  $I = X$  and  $h = a$  in Ring.ideal-subset[of  $Vr\ K\ v$ ],
  assumption+,
  frule-tac  $x = a$  in elem-out-vp-unit[of  $v$ ], assumption+)

apply (frule-tac  $x = a$  and  $I = X$  in ideal-inc-elemOval-whole [of  $v$ ],
  simp add:Vr-mem-f-mem, assumption+)

apply (rule subsetI, simp, erule disjE,
  simp add:prime-ideal-def, simp add:vp-ideal,
  simp add:Ring.whole-ideal, rule subsetI, simp add:vp-mem-Vr-mem)
done

lemma (in Corps) ideal-sub-vp: $\llbracket valuation\ K\ v; ideal\ (Vr\ K\ v)\ I;$ 
 $I \neq carrier\ (Vr\ K\ v) \rrbracket \implies I \subseteq (vp\ K\ v)$ 
apply (frule Vr-ring[of  $v$ ], rule contrapos-pp, simp+)
apply (simp add:subset-eq,
  erule bexE)
apply (frule-tac  $h = x$  in Ring.ideal-subset[of  $Vr\ K\ v\ I$ ], assumption+,
  frule-tac  $x = x$  in elem-out-vp-unit[of  $v$ ], assumption+,
  frule-tac  $x = x$  in ideal-inc-elemOval-whole[of  $v - I$ ],
  simp add:Vr-mem-f-mem, assumption+, simp)
done

lemma (in Corps) Vr-local: $\llbracket valuation\ K\ v; maximal-ideal\ (Vr\ K\ v)\ I \rrbracket \implies$ 
 $(vp\ K\ v) = I$ 
apply (frule Vr-ring[of  $v$ ],
  frule ideal-sub-vp[of  $v\ I$ ], simp add:Ring.maximal-ideal-ideal)
apply (simp add:maximal-ideal-def,
  frule conjunct2, fold maximal-ideal-def, frule conjunct1,
  rule Ring.proper-ideal, assumption+, simp add:maximal-ideal-def, assumption)
apply (rule equalityI) prefer 2 apply assumption
apply (rule contrapos-pp, simp+,
  frule sets-not-eq[of  $vp\ K\ v\ I$ ], assumption+, erule bexE)
apply (frule-tac  $x = a$  in vp-mem-Vr-mem[of  $v$ ],
  frule Ring.maximal-ideal-ideal[of  $Vr\ K\ v\ I$ ], assumption,
  frule-tac  $x = a$  in Ring.elem-out-ideal-nonzero[of  $Vr\ K\ v\ I$ ],
  assumption+,
  frule vp-ideal[of  $v$ ], rule Ring.ideal-subset[of  $Vr\ K\ v\ vp\ K\ v$ ],
  assumption+)

apply (frule-tac  $a = a$  in Ring.principal-ideal[of  $Vr\ K\ v$ ], assumption+,
  frule Ring.maximal-ideal-ideal[of  $Vr\ K\ v\ I$ ], assumption+,
  frule-tac  $?I2.0 = Vr\ K\ v \triangleleft_p a$  ain Ring.sum-ideals[of  $Vr\ K\ v\ I$ ],
  simp add:Ring.maximal-ideal-ideal, assumption,
  frule-tac  $?I2.0 = Vr\ K\ v \triangleleft_p a$  ain Ring.sum-ideals-la1[of  $Vr\ K\ v\ I$ ],
  assumption+,

```

```

frule-tac ?I2.0 = Vr K v  $\Diamond_p$  ain Ring.sum-ideals-la2[of Vr K v I],
  assumption+,
frule-tac a = a in Ring.a-in-principal[of Vr K v], assumption+,
frule-tac A = Vr K v  $\Diamond_p$  a and B = I  $\mp$  (Vr K v) (Vr K v  $\Diamond_p$  a)
  and c = a in subsetD, assumption+)
thm Ring.sum-ideals-cont[of Vr K v vp K v I ]
apply (frule-tac B = Vr K v  $\Diamond_p$  a in Ring.sum-ideals-cont[of Vr K v
  vp K v I], simp add:vp-ideal, assumption)
apply (frule-tac a = a in Ring.ideal-cont-Rxa[of Vr K v vp K v],
  simp add:vp-ideal, assumption+)
apply (simp add:maximal-ideal-def, (erule conjE)+,
  subgoal-tac I  $\mp$  (Vr K v) (Vr K v  $\Diamond_p$  a)  $\in$  {J. ideal (Vr K v) J  $\wedge$  I  $\subseteq$  J},
  simp, thin-tac {J. ideal (Vr K v) J  $\wedge$  I  $\subseteq$  J} = {I, carrier (Vr K v)})
apply (erule disjE, simp)
apply (cut-tac A = carrier (Vr K v) and B = I  $\mp$  Vr K v Vr K v  $\Diamond_p$  a and
  C = vp K v in subset-trans, simp, assumption,
  frule Ring.ideal-subset1[of Vr K v vp K v], simp add:vp-ideal,
  frule equalityI[of vp K v carrier (Vr K v)], assumption+,
  frule vp-not-whole[of v], simp)
apply blast
done

```

```

lemma (in Corps) v-residue-field:valuation K v  $\implies$ 
  Corps ((Vr K v) /r (vp K v))

```

```

by (frule Vr-ring[of v],
  rule Ring.residue-field-cd [of Vr K v vp K v], assumption+,
  simp add:vp-maximal)

```

```

lemma (in Corps) Vr-n-val-Vr:valuation K v  $\implies$ 
  carrier (Vr K v) = carrier (Vr K (n-val K v))
by (simp add:Vr-def Sr-def,
  rule equalityI,
  (rule subsetI, simp, erule conjE, simp add:val-pos-n-val-pos),
  (rule subsetI, simp, erule conjE, simp add:val-pos-n-val-pos[THEN sym]))

```

2.4 Ideals in a valuation ring

```

lemma (in Corps) Vr-has-poss-elem:valuation K v  $\implies$ 
   $\exists x \in \text{carrier (Vr K v)} - \{0_{Vr K v}\}. 0 < v x$ 

```

```

apply (frule val-Pg[of v], erule conjE,
  frule Lv-pos[of v], drule sym,
  subst Vr-0-f-0, assumption+)
apply (frule aeq-ale[of Lv K v v (Pg K v)],
  frule aless-le-trans[of 0 Lv K v v (Pg K v)], assumption+,
  frule val-poss-mem-Vr[of v Pg K v],
  simp, assumption, blast)
done

```


lemma (in Corps) *vp-nonzero:valuation* $K\ v \implies vp\ K\ v \neq \{0_{Vr\ K\ v}\}$
apply (frule *Vr-has-poss-elem*[of v], erule *bexE*,
simp, erule *conjE*,
frule-tac $x1 = x$ in *vp-mem-val-poss*[*THEN sym, of v*],
simp add: *Vr-mem-f-mem*, simp, rule *contrapos-pp*, simp+)
done

lemma (in Corps) *field-frac-mul*: $\llbracket x \in carrier\ K; y \in carrier\ K; y \neq 0 \rrbracket$
 $\implies x = (x \cdot_r (y^{-K})) \cdot_r y$
apply (cut-tac *invf-closed*[of y],
cut-tac *field-is-ring*,
simp add: *Ring.ring-tOp-assoc*,
subst *linvf*[of y], simp, simp add: *Ring.ring-r-one*[of K], simp)
done

lemma (in Corps) *elems-le-val*: $\llbracket valuation\ K\ v; x \in carrier\ K; y \in carrier\ K;$
 $x \neq 0; v\ x \leq (v\ y) \rrbracket \implies \exists r \in carrier\ (Vr\ K\ v). y = r \cdot_r x$
apply (frule *ale-diff-pos*[of $v\ x\ v\ y$], simp add: *diff-ant-def*,
simp add: *value-of-inv*[*THEN sym, of v x*],
cut-tac *invf-closed*[of x],
simp only: *val-t2p*[*THEN sym, of v y x^{-K}*])
apply (cut-tac *field-is-ring*,
frule-tac $x = y$ and $y = x^{-K}$ in *Ring.ring-tOp-closed*[of K],
assumption+,
simp add: *val-pos-mem-Vr*[of $v\ y \cdot_r (x^{-K})$],
frule *field-frac-mul*[of $y\ x$], assumption+, blast)
apply simp
done

lemma (in Corps) *val-Rxa-gt-a*: $\llbracket valuation\ K\ v; x \in carrier\ (Vr\ K\ v) - \{0\};$
 $y \in carrier\ (Vr\ K\ v); y \in Rxa\ (Vr\ K\ v)\ x \rrbracket \implies v\ x \leq (v\ y)$
apply (simp add: *Rxa-def*,
erule *bexE*,
simp add: *Vr-tOp-f-tOp*, (erule *conjE*) +,
frule-tac $x = r$ in *Vr-mem-f-mem*[of v], assumption+,
frule-tac $x = x$ in *Vr-mem-f-mem*[of v], assumption+)
apply (subst *val-t2p*, assumption+,
simp add: *val-pos-mem-Vr*[*THEN sym, of v*],
frule-tac $y = v\ r$ in *aadd-le-mono*[of $0 - v\ x$],
simp add: *aadd-0-l*)
done

lemma (in Corps) *val-Rxa-gt-a-1*: $\llbracket valuation\ K\ v; x \in carrier\ (Vr\ K\ v);$
 $y \in carrier\ (Vr\ K\ v); x \neq 0; v\ x \leq (v\ y) \rrbracket \implies y \in Rxa\ (Vr\ K\ v)\ x$
apply (frule-tac $x = x$ in *Vr-mem-f-mem*[of v], assumption+,
frule-tac $x = y$ in *Vr-mem-f-mem*[of v], assumption+,
frule *v-ale-diff*[of $v\ x\ y$], assumption+,
cut-tac *invf-closed*[of x],
cut-tac *field-is-ring*, frule *Ring.ring-tOp-closed*[of $K\ y\ x^{-K}$],

```

      assumption+)
apply (simp add:val-pos-mem-Vr[of v y ·r (x-K)],
      frule field-frac-mul[of y x], assumption+,
      simp add:Rxa-def, simp add:Vr-tOp-f-tOp, blast, simp)
done

lemma (in Corps) eqval-inv:⟦valuation K v; x ∈ carrier K; y ∈ carrier K;
  y ≠ 0; v x = v y⟧ ⟹ 0 = v (x ·r (y-K))
by (cut-tac invf-closed[of y],
  simp add:val-t2p value-of-inv, simp add:aadd-minus-r,
  simp)

lemma (in Corps) eq-val-eq-idealTr:⟦valuation K v;
  x ∈ carrier (Vr K v) - {0}; y ∈ carrier (Vr K v); v x ≤ (v y)⟧ ⟹
  Rxa (Vr K v) y ⊆ Rxa (Vr K v) x
apply (frule val-Rxa-gt-a-1[of v x y], simp+,
  erule conjE)
apply (frule-tac x = x in Vr-mem-f-mem[of v], assumption+,
  frule Vr-ring[of v],
  frule Ring.principal-ideal[of Vr K v x], assumption,
  frule Ring.ideal-cont-Rxa[of Vr K v (Vr K v) ◇p x y],
  assumption+)
done

lemma (in Corps) eq-val-eq-ideal:⟦valuation K v;
  x ∈ carrier (Vr K v); y ∈ carrier (Vr K v); v x = v y⟧
  ⟹ Rxa (Vr K v) x = Rxa (Vr K v) y
apply (case-tac x = 0K,
  simp add:value-of-zero,
  frule value-inf-zero[of v y],
  simp add:Vr-mem-f-mem, rule sym, assumption, simp)
apply (rule equalityI,
  rule eq-val-eq-idealTr[of v y x], assumption+,
  drule sym, simp,
  rule contrapos-pp, simp+, simp add:value-of-zero,
  frule Vr-mem-f-mem[of v x], assumption+,
  frule value-inf-zero[of v x], assumption+,
  rule sym, assumption, simp, simp, simp)
apply (rule eq-val-eq-idealTr[of v x y], assumption+, simp,
  assumption, rule aeq-ale, assumption+)
done

lemma (in Corps) eq-ideal-eq-val:⟦valuation K v; x ∈ carrier (Vr K v);
  y ∈ carrier (Vr K v); Rxa (Vr K v) x = Rxa (Vr K v) y⟧ ⟹ v x = v y
apply (case-tac x = 0K, simp,
  drule sym,
  frule Vr-ring[of v],
  frule Ring.a-in-principal[of Vr K v y], assumption+, simp,
  thin-tac Vr K v ◇p y = Vr K v ◇p (0), simp add:Rxa-def,

```

```

    erule bexE, simp add: Vr-0-f-0[of v, THEN sym])
  apply (simp add: Vr-tOp-f-tOp, simp add: Vr-0-f-0,
    frule-tac  $x = r$  in Vr-mem-f-mem[of v], assumption+,
    cut-tac field-is-ring, simp add: Ring.ring-times-x-0)
  apply (frule Vr-ring[of v],
    frule val-Rxa-gt-a[of v x y], simp,
    simp)
  apply (drule sym,
    frule Ring.a-in-principal[of Vr K v y], simp, simp)
  apply (frule val-Rxa-gt-a[of v y x],
    simp, rule contrapos-pp, simp+,
    frule Ring.a-in-principal[of Vr K v x], assumption+,
    simp add: Rxa-def,
    erule bexE, simp add: Vr-tOp-f-tOp, cut-tac field-is-ring,
    frule-tac  $x = r$  in Vr-mem-f-mem[of v], assumption+,
    simp add: Ring.ring-times-x-0, simp,
    frule Ring.a-in-principal[of Vr K v x], assumption+, simp,
    rule ale-antisym, assumption+)
done

lemma (in Corps) zero-val-gen-whole:
   $\llbracket \text{valuation } K \ v; \ x \in \text{carrier } (Vr \ K \ v) \rrbracket \implies$ 
     $(v \ x = 0) = (Rxa \ (Vr \ K \ v) \ x = \text{carrier } (Vr \ K \ v))$ 
  apply (frule Vr-mem-f-mem[of v x], assumption,
    frule Vr-ring[of v])
  apply (rule iffI,
    frule Ring.principal-ideal[of Vr K v x], assumption+,
    frule Ring.a-in-principal[of Vr K v x], assumption+,
    rule ideal-inc-elemOval-whole[of v x Vr K v  $\Diamond_p$  x], assumption+,
    frule Ring.ring-one[of Vr K v],
    frule eq-set-inc[of  $1_r(Vr \ K \ v)$ 
      carrier (Vr K v) Vr K v  $\Diamond_p$  x], drule sym, assumption,
    thin-tac  $1_r(Vr \ K \ v) \in \text{carrier } (Vr \ K \ v)$ ,
    thin-tac Vr K v  $\Diamond_p$  x = carrier (Vr K v))
  apply (simp add: Rxa-def, erule bexE,
    simp add: Vr-1-f-1, simp add: Vr-tOp-f-tOp,
    frule value-of-one[of v], simp,
    frule-tac  $x = r$  in Vr-mem-f-mem[of v], assumption+,
    cut-tac field-is-ring, simp add: val-t2p,
    simp add: val-pos-mem-Vr[THEN sym, of v],
    rule contrapos-pp, simp+,
    cut-tac less-le[THEN sym, of 0 v x], drule not-sym, simp,
    frule-tac  $x = v \ r$  and  $y = v \ x$  in aadd-pos-poss, assumption+,
    simp)
done

lemma (in Corps) elem-nonzeroval-gen-proper:  $\llbracket \text{valuation } K \ v; \$ 
   $x \in \text{carrier } (Vr \ K \ v); \ v \ x \neq 0 \rrbracket \implies Rxa \ (Vr \ K \ v) \ x \neq \text{carrier } (Vr \ K \ v)$ 
  apply (rule contrapos-pp, simp+)

```

apply (*simp add: zero-val-gen-whole*[*THEN sym*])
done

We prove that $Vr\ K\ v$ is a principal ideal ring

definition

$LI :: [(r, m)\ Ring-scheme, r \Rightarrow ant, r\ set] \Rightarrow ant$ **where**

$LI\ K\ v\ I = AMin\ (v\ 'I)$

definition

$Ig :: [(r, m)\ Ring-scheme, r \Rightarrow ant, r\ set] \Rightarrow r$ **where**

$Ig\ K\ v\ I = (SOME\ x.\ x \in I \wedge v\ x = LI\ K\ v\ I)$

lemma (**in** *Corps*) *val-in-image*: $\llbracket valuation\ K\ v; ideal\ (Vr\ K\ v)\ I; x \in I \rrbracket \Longrightarrow$
 $v\ x \in v\ 'I$

by (*simp add: image-def, blast*)

lemma (**in** *Corps*) *I-vals-nonempty*: $\llbracket valuation\ K\ v; ideal\ (Vr\ K\ v)\ I \rrbracket \Longrightarrow$
 $v\ 'I \neq \{\}$

by (*frule Vr-ring*[*of v*],
frule Ring.ideal-zero[*of Vr K v I*],
assumption+, *rule contrapos-pp*, *simp+*)

lemma (**in** *Corps*) *I-vals-LBset*: $\llbracket valuation\ K\ v; ideal\ (Vr\ K\ v)\ I \rrbracket \Longrightarrow$
 $v\ 'I \subseteq LBset\ 0$

apply (*frule Vr-ring*[*of v*],
rule subsetI, *simp add: LBset-def*, *simp add: image-def*)

apply (*erule bexE*,
frule-tac h = xa in Ring.ideal-subset[*of Vr K v I*], *assumption+*)

apply (*frule-tac x1 = xa in val-pos-mem-Vr*[*THEN sym, of v*],
simp add: Vr-mem-f-mem, *simp*)

done

lemma (**in** *Corps*) *LI-pos*: $\llbracket valuation\ K\ v; ideal\ (Vr\ K\ v)\ I \rrbracket \Longrightarrow 0 \leq LI\ K\ v\ I$

apply (*simp add: LI-def*,
frule I-vals-LBset[*of v*],
simp add: ant-0[*THEN sym*],
frule I-vals-nonempty[*of v*], *simp only: ant-0*)

apply (*simp only: ant-0*[*THEN sym*], *frule AMin*[*of v 'I 0*], *assumption*,
erule conjE, *frule subsetD*[*of v 'I LBset (ant 0) AMin (v 'I)*],
assumption+, *simp add: LBset-def*)

done

lemma (**in** *Corps*) *LI-poss*: $\llbracket valuation\ K\ v; ideal\ (Vr\ K\ v)\ I; I \neq carrier\ (Vr\ K\ v) \rrbracket \Longrightarrow 0 < LI\ K\ v\ I$

apply (*subst less-le*)
apply (*simp add: LI-pos*)

apply (*rule contrapos-pp, simp+*)

apply (*simp add:LI-def,*
frule I-vals-LBset[of v], assumption+,
simp add:ant-0[THEN sym],
frule I-vals-nonempty[of v], assumption+, simp only:ant-0)

apply (*simp only:ant-0[THEN sym], frule AMin[of v ' I 0], assumption,*
erule conjE, frule subsetD[of v ' I LBset (ant 0) AMin (v ' I)],
assumption+, simp add:LBset-def)

apply (*thin-tac $\forall x \in I. \text{ant } 0 \leq v x,$*
thin-tac $v ' I \subseteq \{x. \text{ant } 0 \leq x\},$ simp add:image-def,
erule bexE, simp add:ant-0)

apply (*frule Vr-ring[of v],*
frule-tac $h = x$ in Ring.ideal-subset[of Vr K v I], assumption+,
frule-tac $x = x$ in zero-val-gen-whole[of v], assumption+,
simp,
frule-tac $a = x$ in Ring.ideal-cont-Rxa[of Vr K v I], assumption+,
simp, frule Ring.ideal-subset1[of Vr K v I], assumption+)

apply (*frule equalityI[of I carrier (Vr K v)], assumption+, simp*)
done

lemma (*in Corps*) *LI-z: [valuation K v; ideal (Vr K v) I; $I \neq \{0_{Vr K v}\}] \implies$*
 $\exists z. LI K v I = \text{ant } z$

apply (*frule Vr-ring[of v],*
frule Ring.ideal-zero[of Vr K v I], assumption+,
cut-tac mem-ant[of LI K v I],
frule LI-pos[of v I], assumption,
erule disjE, simp,
cut-tac minf-le-any[of 0],
frule ale-antisym[of 0 $-\infty$], assumption+, simp)

apply (*erule disjE, simp,*
frule singleton-sub[of $0_{Vr K v} I$],
frule sets-not-eq[of I $\{0_{Vr K v}\}$], assumption+,
erule bexE, simp)

apply (*simp add:LI-def,*
frule I-vals-LBset[of v], assumption+,
simp only:ant-0[THEN sym],
frule I-vals-nonempty[of v], assumption+,
frule AMin[of v ' I 0], assumption, erule conjE)

apply (*frule-tac $x = a$ in val-in-image[of v I], assumption+,*
drule-tac $x = v a$ in bspec, simp,
simp add:Vr-0-f-0,
frule-tac $x = a$ in val-nonzero-z[of v],
simp add:Ring.ideal-subset Vr-mem-f-mem, assumption+,
erule exE, simp,
cut-tac $x = \text{ant } z$ in inf-ge-any, frule-tac $x = \text{ant } z$ in

$\text{ale-antisym}[\text{of } - \infty], \text{assumption+}, \text{simp})$
done

lemma (**in** *Corps*) *LI-k*: $\llbracket \text{valuation } K \text{ } v; \text{ideal } (Vr \text{ } K \text{ } v) \text{ } I \rrbracket \implies$
 $\exists k \in I. LI \text{ } K \text{ } v \text{ } I = v \text{ } k$

by (*simp add:LI-def*,
frule I-vals-LBset[*of v*], *assumption+*,
simp only:ant-0[*THEN sym*],
frule I-vals-nonempty[*of v*], *assumption+*,
frule AMin[*of v ' I 0*], *assumption*, *erule conjE*,
thin-tac $\forall x \in v \text{ ' } I. AMin \text{ } (v \text{ ' } I) \leq x$, *simp add:image-def*)

lemma (**in** *Corps*) *LI-infinity*: $\llbracket \text{valuation } K \text{ } v; \text{ideal } (Vr \text{ } K \text{ } v) \text{ } I \rrbracket \implies$
 $(LI \text{ } K \text{ } v \text{ } I = \infty) = (I = \{\mathbf{0}_{Vr \text{ } K \text{ } v}\})$

apply (*frule Vr-ring*[*of v*])
apply (*rule iffI*)
apply (*rule contrapos-pp*, *simp+*,
frule Ring.ideal-zero[*of Vr K v I*], *assumption+*,
frule singleton-sub[*of 0_{Vr K v I}*],
frule sets-not-eq[*of I {0_{Vr K v}}*], *assumption+*,
erule bexE,
frule-tac $h = a \text{ in } Ring.ideal-subset[\text{of } Vr \text{ } K \text{ } v \text{ } I]$, *assumption+*,
simp add:Vr-0-f-0,
frule-tac $x = a \text{ in } Vr-mem-f-mem[\text{of } v]$, *assumption+*,
frule-tac $x = a \text{ in } val-nonzero-z[\text{of } v]$, *assumption+*,
erule exE,
simp add:LI-def,
frule I-vals-LBset[*of v*], *assumption+*,
simp only:ant-0[*THEN sym*],
frule I-vals-nonempty[*of v*], *assumption+*,
frule AMin[*of v ' I 0*], *assumption*, *erule conjE*)
apply (*frule-tac* $h = a \text{ in } Ring.ideal-subset[\text{of } Vr \text{ } K \text{ } v \text{ } I]$, *assumption+*,
frule-tac $x = a \text{ in } val-in-image[\text{of } v \text{ } I]$, *assumption+*,
drule-tac $x = v \text{ } a \text{ in } bspec$, *simp*)
apply (*frule-tac* $x = a \text{ in } val-nonzero-z[\text{of } v]$, *assumption+*,
erule exE, *simp*,
cut-tac $x = ant \text{ } z \text{ in } inf-ge-any$, *frule-tac* $x = ant \text{ } z \text{ in}$
 $\text{ale-antisym}[\text{of } - \infty], \text{assumption+}, \text{simp})$

apply (*frule sym*, *thin-tac* $I = \{\mathbf{0}_{Vr \text{ } K \text{ } v}\}$,
simp add:LI-def,
frule I-vals-LBset[*of v*], *assumption+*,
simp only:ant-0[*THEN sym*],
frule I-vals-nonempty[*of v*], *assumption+*,
frule AMin[*of v ' I 0*], *assumption*, *erule conjE*,
drule sym, *simp*,
simp add:Vr-0-f-0 value-of-zero)
done

lemma (in Corps) *val-Ig*: $\llbracket \text{valuation } K \ v; \text{ideal } (Vr \ K \ v) \ I \rrbracket \implies$
 $(Ig \ K \ v \ I) \in I \wedge v \ (Ig \ K \ v \ I) = LI \ K \ v \ I$

by (*simp add*:*Ig-def*, *rule someI2-ex*,
frule LI-k[of *v I*], *assumption+*, *erule bexE*,
drule sym, *blast*, *assumption*)

lemma (in Corps) *Ig-nonzero*: $\llbracket \text{valuation } K \ v; \text{ideal } (Vr \ K \ v) \ I; I \neq \{\mathbf{0}_{Vr \ K \ v}\} \rrbracket$
 $\implies (Ig \ K \ v \ I) \neq \mathbf{0}$

by (*rule contrapos-pp*, *simp+*,
frule LI-infinity[of *v I*], *assumption+*,
frule val-Ig[of *v I*], *assumption+*, *erule conjE*,
simp add:*value-of-zero*)

lemma (in Corps) *Vr-ideal-npowf-closed*: $\llbracket \text{valuation } K \ v; \text{ideal } (Vr \ K \ v) \ I;$
 $x \in I; 0 < n \rrbracket \implies x_K^n \in I$

by (*simp add*:*npowf-def*, *frule Vr-ring*[of *v*],
frule Ring.ideal-mpow-closed[of *Vr K v I x nat n*], *assumption+*,
simp, *frule Ring.ideal-subset*[of *Vr K v I x*], *assumption+*,
simp add:*Vr-exp-f-exp*)

lemma (in Corps) *Ig-generate-I*: $\llbracket \text{valuation } K \ v; \text{ideal } (Vr \ K \ v) \ I \rrbracket \implies$
 $(Vr \ K \ v) \diamond_p (Ig \ K \ v \ I) = I$

apply (*frule Vr-ring*[of *v*])

apply (*case-tac* $I = \text{carrier } (Vr \ K \ v)$,
frule sym, *thin-tac* $I = \text{carrier } (Vr \ K \ v)$,
frule Ring.ring-one[of *Vr K v*],
simp, *simp add*: *Vr-1-f-1*,
frule val-Ig[of *v I*], *assumption+*, *erule conjE*,
frule LI-pos[of *v I*], *assumption+*,

simp add: *LI-def cong del: image-cong-simp*,
frule I-vals-LBset[of *v*], *assumption+*,
simp only: *ant-0*[*THEN sym*],
frule I-vals-nonempty[of *v*], *assumption+*,
frule AMin[of *v ' I 0*], *assumption*, *erule conjE*,

frule val-in-image[of *v I 1_r*], *assumption+*,
drule-tac $x = v \ (1_r)$ **in** *bspec*, *assumption+*,
simp add: *value-of-one ant-0 cong del: image-cong-simp*,
simp add: *zero-val-gen-whole*[of *v Ig K v I*])

apply (*frule val-Ig*[of *v I*], *assumption+*, (*erule conjE*) $+$,
frule Ring.ideal-cont-Rxa[of *Vr K v I Ig K v I*], *assumption+*,
rule equalityI, *assumption+*)

apply (*case-tac* $LI \ K \ v \ I = \infty$,
frule LI-infinity[of *v I*], *simp*,
simp add:*Rxa-def*, *simp add*:*Ring.ring-times-x-0*,
frule Ring.ring-zero, *blast*)

apply (*rule subsetI*,
case-tac $v\ x = 0$,
frule-tac $x = x$ **in** *Vr-mem-f-mem*[*of v*],
simp add:Ring.ideal-subset,
frule-tac $x = x$ **in** *zero-val-gen-whole*[*of v*],
simp add:Ring.ideal-subset, *simp*,
frule-tac $a = x$ **in** *Ring.ideal-cont-Rxa*[*of Vr K v I*], *assumption+*,
simp, *frule Ring.ideal-subset1*[*of Vr K v I*], *assumption*,
frule equalityI[*of I carrier (Vr K v)*], *assumption+*, *simp*)
apply (*simp add:LI-def*,
frule I-vals-LBset[*of v*], *assumption+*,
simp only:ant-0[THEN sym],
frule I-vals-nonempty[*of v*], *assumption+*,
frule AMin[*of v ' I 0*], *assumption*, *erule conjE*,
frule-tac $x = v\ x$ **in** *bspec*,
frule-tac $x = x$ **in** *val-in-image*[*of v I*], *assumption+*,
simp)
apply (*drule-tac* $x = x$ **in** *bspec*, *assumption*,
frule-tac $y = x$ **in** *eq-val-eq-idealTr*[*of v Ig K v I*],
simp add:Ring.ideal-subset,
rule contrapos-pp, *simp+*, *simp add:value-of-zero*,
simp add:Ring.ideal-subset, *simp*)

apply (*frule-tac* $a = x$ **in** *Ring.a-in-principal*[*of Vr K v*],
simp add:Ring.ideal-subset, *rule subsetD*, *assumption+*)
done

lemma (**in** *Corps*) *Pg-gen-vp:valuation K v* $\implies$
 $(Vr\ K\ v) \triangleleft_p (Pg\ K\ v) = vp\ K\ v$
apply (*frule vp-ideal*[*of v*],
frule Ig-generate-I[*of v vp K v*], *assumption+*,
frule vp-not-whole[*of v*],
frule eq-val-eq-ideal[*of v Ig K v (vp K v) Pg K v*],
frule val-Ig [*of v vp K v*], *assumption+*, *erule conjE*,
simp add:vp-mem-Vr-mem)

apply (*frule val-Pg*[*of v*], *erule conjE*,
frule Lv-pos[*of v*],
rotate-tac -2, *drule sym*, *simp*,
simp add:val-poss-mem-Vr)

apply (*thin-tac* $Vr\ K\ v \triangleleft_p Ig\ K\ v (vp\ K\ v) = vp\ K\ v$,
frule val-Pg[*of v*], *erule conjE*,
simp, *frule val-Ig*[*of v vp K v*], *assumption+*, *erule conjE*,
simp, *thin-tac* $v (Pg\ K\ v) = Lv\ K\ v$,
thin-tac $Ig\ K\ v (vp\ K\ v) \in vp\ K\ v \wedge v (Ig\ K\ v (vp\ K\ v)) =$
 $LI\ K\ v (vp\ K\ v)$, *simp add:LI-def Lv-def*,
subgoal-tac $v \triangleleft vp\ K\ v = \{x. x \in v \triangleleft carrier\ K \wedge 0 < x\}$,

$\text{simp})$

apply ($\text{thin-tac ideal (Vr K v) (vp K v), thin-tac Pg K v} \in \text{carrier K,}$
 $\text{thin-tac Pg K v} \neq \mathbf{0},$
 $\text{rule equalityI, rule subsetI,}$
 $\text{simp add:image-def vp-def, erule exE, erule conjE,}$
 $(\text{erule conjE})+,$
 $\text{frule-tac } x = xa \text{ in Vr-mem-f-mem[of v], assumption+, simp, blast})$

apply ($\text{rule subsetI, simp add:image-def vp-def, erule conjE, erule bexE, simp,}$
 $\text{frule-tac } x = xa \text{ in val-poss-mem-Vr[of v], assumption+,}$
 $\text{cut-tac } y = v \text{ xa in less-le[of 0], simp, blast, simp})$

done

lemma (**in** Corps) $\text{vp-gen-t:valuation K v} \implies$
 $\exists t \in \text{carrier (Vr K v). vp K v} = (\text{Vr K v}) \diamond_p t$

by ($\text{frule Pg-gen-vp[of v], frule Pg-in-Vr[of v], blast})$

lemma (**in** Corps) $\text{vp-gen-nonzero:} \llbracket \text{valuation K v; vp K v} = (\text{Vr K v}) \diamond_p t \rrbracket \implies$
 $t \neq \mathbf{0}_{\text{Vr K v}}$

apply ($\text{rule contrapos-pp, simp+,}$
 $\text{cut-tac Ring.Rxa-zero[of Vr K v], drule sym, simp,}$
 $\text{simp add:vp-nonzero})$

apply ($\text{simp add:Vr-ring})$

done

lemma (**in** Corps) $\text{n-value-idealTr:} \llbracket \text{valuation K v; } 0 \leq n \rrbracket \implies$
 $(\text{vp K v}) \diamond (\text{Vr K v})^n = \text{Vr K v} \diamond_p ((\text{Pg K v})^{\wedge (\text{Vr K v})^n})$

apply ($\text{frule Vr-ring[of v],}$
 $\text{frule Pg-gen-vp[THEN sym, of v],}$
 $\text{simp add:vp-ideal,}$
 $\text{frule val-Pg[of v], simp, (erule conjE)+})$

apply ($\text{subst Ring.principal-ideal-n-pow[of Vr K v Pg K v}$
 $\text{Vr K v} \diamond_p \text{Pg K v], assumption+,}$
 $\text{frule Lv-pos[of v], rotate-tac -2, frule sym,}$
 $\text{thin-tac } v (\text{Pg K v}) = \text{Lv K v, simp, simp add:val-poss-mem-Vr,}$
 $\text{simp+})$

done

lemma (**in** Corps) $\text{ideal-pow-vp:} \llbracket \text{valuation K v; ideal (Vr K v) I;}$
 $\text{I} \neq \text{carrier (Vr K v); I} \neq \{\mathbf{0}_{\text{Vr K v}}\} \rrbracket \implies$
 $\text{I} = (\text{vp K v}) \diamond (\text{Vr K v}) (\text{na (n-val K v (Ig K v I))})$

apply ($\text{frule Vr-ring[of v],}$
 $\text{frule Ig-generate-I[of v I], assumption+})$

apply ($\text{frule n-val[of v Ig K v I],}$
 $\text{frule val-Ig[of v I], assumption+, erule conjE,}$
 $\text{simp add:Ring.ideal-subset[of Vr K v I Ig K v I] Vr-mem-f-mem})$

apply (frule val-Pg[of v], erule conjE,
 rotate-tac -1, drule sym, simp,
 frule Ig-nonzero[of v I], assumption+,
 frule LI-pos[of v I], assumption+,
 frule Lv-pos[of v],
 frule val-Ig[of v I], assumption+, (erule conjE)+,
 rotate-tac -1, drule sym, simp,
 frule val-pos-n-val-pos[of v Ig K v I],
 simp add:Ring.ideal-subset Vr-mem-f-mem,
 simp)
apply (frule zero-val-gen-whole[THEN sym, of v Ig K v I],
 simp add:Ring.ideal-subset,
 simp, rotate-tac -1, drule not-sym,
 cut-tac less-le[THEN sym, of 0 v (Ig K v I)], simp,
 thin-tac $0 \leq v$ (Ig K v I),
 frule Ring.ideal-subset[of Vr K v I Ig K v I], assumption+,
 frule Vr-mem-f-mem[of v Ig K v I], assumption+,
 frule val-poss-n-val-poss[of v Ig K v I], assumption+, simp)
apply (frule Ig-nonzero[of v I],
 frule val-nonzero-noninf[of v Ig K v I], assumption+,
 simp add:val-noninf-n-val-noninf[of v Ig K v I],
 frule val-poss-mem-Vr[of v Pg K v], assumption+,
 subst n-value-idealTr[of v na (n-val K v (Ig K v I))],
 assumption+, simp add:na-def)

apply (frule eq-val-eq-ideal[of v Ig K v I
 (Pg K v) $\neg$ (Vr K v) (na (n-val K v (Ig K v I)))], assumption+,
 rule Ring.npClose, assumption+,
 simp add:Vr-exp-f-exp[of v Pg K v],
 subst val-exp-ring[THEN sym, of v Pg K v
 na (n-val K v (Ig K v I))], assumption+)
apply (frule Lv-z[of v], erule exE, simp,
 rotate-tac 6, drule sym, simp,
 subst asprod-amult,
 simp add:val-poss-n-val-poss[of v Ig K v I],
 frule val-nonzero-noninf[of v Ig K v I], assumption+,
 frule val-noninf-n-val-noninf[of v Ig K v I], assumption+, simp,
 rule aposs-na-poss[of n-val K v (Ig K v I)], assumption+)
apply (fold an-def)
apply (subst an-na[THEN sym, of n-val K v (Ig K v I)],
 frule val-nonzero-noninf[of v Ig K v I], assumption+,
 frule val-noninf-n-val-noninf[of v Ig K v I], assumption+, simp,
 simp add:aless-imp-le, simp)
apply simp
done

lemma (in Corps) ideal-apow-vp: [valuation K v; ideal (Vr K v) I] $\implies$
 $I = (vp K v) (Vr K v) (n-val K v (Ig K v I))$
apply (frule Vr-ring[of v])

apply (*case-tac* v ($Ig\ K\ v\ I$) = ∞ ,
frule $val-Ig$ [*of* $v\ I$], *assumption*,
frule $val-inf-n-val-inf$ [*of* $v\ Ig\ K\ v\ I$],
simp *add:Ring.ideal-subset* *Vr-mem-f-mem*, *simp*, *simp* *add:r-apow-def*,
simp *add:LI-infinity*[*of* $v\ I$])

apply (*case-tac* v ($Ig\ K\ v\ I$) = 0 ,
frule $val-0-n-val-0$ [*of* $v\ Ig\ K\ v\ I$],
frule $val-Ig$ [*of* $v\ I$], *assumption+*, *erule* *conjE*,
simp *add:Ring.ideal-subset* *Vr-mem-f-mem*, *simp*,

frule $val-Ig$ [*of* $v\ I$], *assumption*,
frule $zero-val-gen-whole$ [*of* $v\ Ig\ K\ v\ I$],
simp *add:Ring.ideal-subset*, (*erule* *conjE*)*+*, *simp*,
frule $Ring.ideal-cont-Rxa$ [*of* $Vr\ K\ v\ I\ Ig\ K\ v\ I$], *assumption+*)

apply (*simp*,
frule $Ring.ideal-subset1$ [*of* $Vr\ K\ v\ I$], *assumption+*,
frule $equalityI$ [*of* $I\ carrier\ (Vr\ K\ v)$], *assumption+*,
simp *add:r-apow-def*)

apply (*frule* $val-noninf-n-val-noninf$ [*of* $v\ Ig\ K\ v\ I$],
frule $val-Ig$ [*of* $v\ I$], *assumption*,
simp *add:Ring.ideal-subset* *Vr-mem-f-mem*, *simp*,
frule $value-n0-n-val-n0$ [*of* $v\ Ig\ K\ v\ I$],
frule $val-Ig$ [*of* $v\ I$], *assumption*,
simp *add:Ring.ideal-subset* *Vr-mem-f-mem*, *assumption*)

apply (*simp* *add:r-apow-def*,
rule $ideal-pow-vp$, *assumption+*,
frule $elem-nonzeroval-gen-proper$ [*of* $v\ Ig\ K\ v\ I$],
frule $val-Ig$ [*of* $v\ I$], *assumption+*, *erule* *conjE*,
simp *add:Ring.ideal-subset*, *assumption*, *simp* *add:Ig-generate-I*)

apply (*frule* $val-Ig$ [*of* $v\ I$], *assumption+*, *erule* *conjE*, *simp*,
simp *add:LI-infinity*[*of* $v\ I$])

done

lemma (*in* *Corps*) *ideal-apow-n-val*: $\llbracket valuation\ K\ v; x \in carrier\ (Vr\ K\ v) \rrbracket \implies$
 $(Vr\ K\ v) \diamond_p x = (vp\ K\ v)(Vr\ K\ v)\ (n-val\ K\ v\ x)$

apply (*frule* $Vr-ring$ [*of* v],
frule $Ring.principal-ideal$ [*of* $Vr\ K\ v\ x$], *assumption+*,
frule $ideal-apow-vp$ [*of* $v\ Vr\ K\ v\ \diamond_p\ x$], *assumption+*)

apply (*frule* $val-Ig$ [*of* $v\ Vr\ K\ v\ \diamond_p\ x$], *assumption+*, *erule* *conjE*,
frule $Ring.ideal-subset$ [*of* $Vr\ K\ v\ Vr\ K\ v\ \diamond_p\ x$
 $Ig\ K\ v\ (Vr\ K\ v\ \diamond_p\ x)$], *assumption+*,
frule $Ig-generate-I$ [*of* $v\ Vr\ K\ v\ \diamond_p\ x$], *assumption+*)

apply (*frule* $eq-ideal-eq-val$ [*of* $v\ Ig\ K\ v\ (Vr\ K\ v\ \diamond_p\ x)$],
assumption+,

$\text{thin-tac } Vr K v \Diamond_p Ig K v (Vr K v \Diamond_p x) = Vr K v \Diamond_p x,$
 $\text{thin-tac } v (Ig K v (Vr K v \Diamond_p x)) = LI K v (Vr K v \Diamond_p x),$
 $\text{frule } n\text{-val}[THEN \text{ sym, of } v x],$
 $\text{simp add: Vr-mem-f-mem, simp,}$
 $\text{thin-tac } v x = n\text{-val } K v x * Lv K v,$
 $\text{frule } n\text{-val}[THEN \text{ sym, of } v Ig K v (Vr K v \Diamond_p x)],$
 $\text{simp add: Vr-mem-f-mem, simp,}$
 $\text{thin-tac } v (Ig K v (Vr K v \Diamond_p x)) = n\text{-val } K v x * Lv K v$
apply ($\text{frule } Lv\text{-pos}[of v],$
 $\text{frule } Lv\text{-z}[of v], \text{erule } exE, \text{simp,}$
 $\text{frule-tac } s = z \text{ in } z\text{less-neq}[THEN \text{ not-sym, of } 0],$
 $\text{frule-tac } z = z \text{ in } a\text{div-eq}[of - n\text{-val } K v (Ig K v (Vr K v \Diamond_p x))$
 $n\text{-val } K v x], \text{assumption+}, \text{simp})$
done

lemma (**in** *Corps*) $t\text{-gen-vp}:\llbracket \text{valuation } K v; t \in \text{carrier } K; v t = 1 \rrbracket \implies$
 $(Vr K v) \Diamond_p t = vp K v$

proof –

assume $a1:\text{valuation } K v$ **and**
 $a2:t \in \text{carrier } K$ **and**
 $a3:v t = 1$
from $a1$ **and** $a2$ **and** $a3$ **have** $h1:t \in \text{carrier } (Vr K v)$
apply ($\text{cut-tac } a0\text{-less-1}$)
apply ($\text{rule } val\text{-poss-mem-}Vr[of v t], \text{assumption+}, \text{simp}$) **done**
from $a1$ **and** $a2$ **and** $a3$ **have** $h2:n\text{-val } K v = v$
apply ($\text{subst } val\text{-surj-n-val}[of v]$) **apply** assumption
apply blast **apply** simp **done**
from $a1$ **and** $h1$ **have** $h3:Vr K v \Diamond_p t = vp K v (Vr K v) (n\text{-val } K v t)$
apply ($\text{simp add:ideal-apow-n-val}[of v t]$) **done**
from $a1$ **and** $a3$ **and** $h2$ **and** $h3$ **show** $?thesis$
apply ($\text{simp add:r-apow-def}$)
apply ($\text{simp only:ant-1}[THEN \text{ sym}], \text{simp only:ant-0}[THEN \text{ sym}]$)
apply ($\text{simp only:aeq-zeq, simp}$)
apply ($\text{cut-tac } z\text{-neg-inf}[THEN \text{ not-sym, of } 1], \text{simp}$)
apply ($\text{simp only:an-1}[THEN \text{ sym}]$) **apply** (simp add:na-an)
apply ($\text{rule } Ring.\text{idealprod-whole-r}[of Vr K v vp K v]$)
apply (simp add:Vr-ring)
apply (simp add:vp-ideal) **done**
qed

lemma (**in** *Corps*) $t\text{-vp-apow}:\llbracket \text{valuation } K v; t \in \text{carrier } K; v t = 1 \rrbracket \implies$
 $(Vr K v) \Diamond_p (t \neg (Vr K v) n) = (vp K v) (Vr K v) (an n)$

proof –

assume $a1:\text{valuation } K v$ **and**
 $a2:t \in \text{carrier } K$ **and**

```

    a3:v t = 1
from a1 have h1:Ring (Vr K v) by (simp add:Vr-ring[of v])
from a1 and a2 and a3 have h2:t ∈ carrier (Vr K v)
  apply (cut-tac a0-less-1)
  apply (rule val-poss-mem-Vr) apply assumption+ apply simp done
from a1 and a2 and a3 and h1 and h2 show ?thesis
  apply (subst Ring.principal-ideal-n-pow[THEN sym, of Vr K v t vp K v n])
  apply assumption+
  apply (simp add:t-gen-vp)
  apply (simp add:r-apow-def)
  apply (rule conjI, rule impI,
    simp only:an-0[THEN sym], frule an-inj[of n 0], simp)
  apply (rule impI)
  apply (rule conjI, rule impI)
  apply (simp add:an-def)
  apply (rule impI, cut-tac an-nat-pos[of n], simp add:na-an)
done
qed

```

```

lemma (in Corps) nonzeroelem-gen-nonzero: [valuation K v; x ≠ 0;
  x ∈ carrier (Vr K v)] ⇒ Vr K v ◇p x ≠ {0Vr K v}
by (frule Vr-ring[of v],
  frule-tac a = x in Ring.a-in-principal[of Vr K v], assumption+,
  rule contrapos-pp, simp+, simp add:Vr-0-f-0)

```

2.4.1 Amin lemma (in Corps)s

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lemma (in Corps) Amin-le-addTr:valuation K v ⇒
(∀ j ≤ n. f j ∈ carrier K) → Amin n (v ∘ f) ≤ (v (nsum K f n))
  apply (induct-tac n)
  apply (rule impI, simp)

  apply (rule impI,
    simp,
    frule-tac x = Σe K f n and y = f (Suc n) in amin-le-plus[of v],
    cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    cut-tac n = n in aGroup.nsum-mem[of K - f], assumption,
    rule allI, simp add:funcset-mem, assumption, simp)
  apply (frule-tac z = Amin n (λu. v (f u)) and z' = v (Σe K f n) and
    w = v (f (Suc n)) in amin-aminTr,
    rule-tac i = amin (Amin n (λu. v (f u))) (v (f (Suc n))) and
    j = amin (v (Σe K f n)) (v (f (Suc n))) and
    k = v (Σe K f n ± (f (Suc n))) in ale-trans, assumption+)
done

```

```

lemma (in Corps) Amin-le-add: [valuation K v; ∀ j ≤ n. f j ∈ carrier K] ⇒
  Amin n (v ∘ f) ≤ (v (nsum K f n))
by (frule Amin-le-addTr[of v n f], simp)

```

lemma (in Corps) value-ge-add: $\llbracket \text{valuation } K \ v; \forall j \leq n. f \ j \in \text{carrier } K; \forall j \leq n. z \leq ((v \circ f) \ j) \rrbracket \implies z \leq (v \ (\Sigma_e \ K \ f \ n))$
apply (frule Amin-le-add[of v n f], assumption+,
cut-tac Amin-ge[of n v $\circ$ f z],
rule ale-trans, assumption+)
apply (rule allI, rule impI,
simp add:comp-def Zset-def,
rule value-in-aug-inf[of v], assumption+, simp+)
done

lemma (in Corps) Vr-ideal-powTr1: $\llbracket \text{valuation } K \ v; \text{ideal } (Vr \ K \ v) \ I; I \neq \text{carrier } (Vr \ K \ v); b \in I \rrbracket \implies b \in (vp \ K \ v)$
by (frule ideal-sub-vp[of v I], assumption+, simp add:subsetD)

2.5 pow of vp and n -value – convergence –

lemma (in Corps) n-value-x-1: $\llbracket \text{valuation } K \ v; 0 \leq n; x \in (vp \ K \ v) \ (Vr \ K \ v) \ n \rrbracket \implies n \leq (n\text{-val } K \ v \ x)$

apply ((case-tac n = ∞ , simp add:r-apow-def,
simp add:Vr-0-f-0, cut-tac field-is-ring,
frule Ring.ring-zero[of K], frule val-inf-n-val-inf[of v 0],
assumption+, simp add:value-of-zero),
(case-tac n = 0, simp add:r-apow-def,
subst val-pos-n-val-pos[THEN sym, of v x], assumption+,
simp add:Vr-mem-f-mem,
subst val-pos-mem-Vr[of v x], assumption+,
simp add:Vr-mem-f-mem, assumption,
simp add:r-apow-def, frule Vr-ring[of v],
frule vp-pow-ideal[of v na n],
frule Ring.ideal-subset[of Vr K v (vp K v) $\Diamond$ (Vr K v) (na n) x],
assumption+, frule Vr-mem-f-mem[of v x], assumption+))

apply (case-tac x = 0_K , simp,
frule value-of-zero[of v],
simp add:val-inf-n-val-inf,
simp add:n-value-idealTr[of v na n],

frule val-Pg[of v], erule conjE, simp, erule conjE,
frule Lv-pos[of v],
rotate-tac -4, frule sym, thin-tac v (Pg K v) = Lv K v, simp,
frule val-poss-mem-Vr[of v Pg K v], assumption+,
frule val-Rxa-gt-a[of v Pg K v $\neg$ (Vr K v) (na n) x],

frule Vr-integral[of v],

```

    simp only: Vr-0-f-0[of v, THEN sym],
    frule Idomain.idom-potent-nonzero[of Vr K v Pg K v na n],
    assumption+, simp, simp add: Ring.npClose, assumption+)

apply (thin-tac  $x \in Vr K v \Diamond_p (Pg K v \neg (Vr K v) (na n))$ ,
    thin-tac ideal (Vr K v) (Vr K v  $\Diamond_p (Pg K v \neg (Vr K v) (na n))$ )))

apply (simp add: Vr-exp-f-exp[of v Pg K v],
    simp add: val-exp-ring[THEN sym, of v Pg K v],
    simp add: n-val[THEN sym, of v x],
    frule val-nonzero-z[of v Pg K v], assumption+,
    erule exE, simp,
    frule aposs-na-poss[of n], simp add: less-le,
    simp add: asprod-amult,

    frule-tac  $w = z$  in amult-pos-mono-r[of - ant (int (na n))
    n-val K v x], simp,
    cut-tac an-na[of n], simp add: an-def, assumption+)
done

lemma (in Corps) n-value-x-1-nat:  $\llbracket \text{valuation } K v; x \in (vp K v) \Diamond (Vr K v) n \rrbracket \implies$ 
     $(an n) \leq (n\text{-val } K v x)$ 
apply (cut-tac an-nat-pos[of n])
apply (frule n-value-x-1[of v an n x], assumption+)
apply (simp add: r-apow-def)
apply (case-tac  $n = 0$ , simp, simp)
apply (cut-tac aless-nat-less[THEN sym, of 0 n])
apply simp
unfolding less-le
apply simp
apply (cut-tac an-neq-inf [of n])
apply simp
apply (simp add: na-an)
apply assumption
done

lemma (in Corps) n-value-x-2:  $\llbracket \text{valuation } K v; x \in \text{carrier } (Vr K v);$ 
     $n \leq (n\text{-val } K v x); 0 \leq n \rrbracket \implies x \in (vp K v) (Vr K v) n$ 
apply (frule Vr-ring[of v],
    frule val-Pg[of v], erule conjE,
    simp, erule conjE, drule sym,
    frule Lv-pos[of v], simp,
    frule val-poss-mem-Vr[of v Pg K v], assumption+)

apply (case-tac  $n = \infty$ ,
    simp add: r-apow-def, cut-tac inf-ge-any[of n-val K v x],
    frule ale-antisym[of n-val K v x  $\infty$ ], assumption+,
    frule val-inf-n-val-inf[THEN sym, of v x],
    simp add: Vr-mem-f-mem, simp,

```

$\text{frule value-inf-zero[of } v \ x],$
 $\text{simp add: Vr-mem-f-mem, simp+}, \text{ simp add: Vr-0-f-0})$

apply ($\text{case-tac } n = 0,$
 $\text{simp add:r-apow-def},$
 $\text{simp add:r-apow-def},$
 $\text{subst n-value-idealTr[of } v \ na \ n], \text{ assumption+},$
 $\text{simp add:apos-na-pos})$

apply ($\text{rule val-Rxa-gt-a-1[of } v \ Pg \ K \ v \neg (Vr \ K \ v) \ (na \ n) \ x],$
 $\text{assumption+},$
 $\text{rule Ring.npClose, assumption+},$
 $\text{simp add: Vr-0-f-0[THEN sym, of } v],$
 $\text{frule Vr-integral[of } v],$
 $\text{frule val-poss-mem-Vr[of } v \ Pg \ K \ v], \text{ assumption+},$
 $\text{simp add:Idomain.idom-potent-nonzero})$

apply ($\text{simp add: Vr-exp-f-exp},$
 $\text{simp add:val-exp-ring[THEN sym, of } v],$
 $\text{rotate-tac } -5, \text{ drule sym},$
 $\text{frule Lv-z[of } v], \text{ erule exE, simp},$
 $\text{frule apos-na-poss[of } n], \text{ simp add: less-le},$
 $\text{simp add:asprod-amult, subst n-val[THEN sym, of } v \ x],$
 $\text{assumption+},$
 $\text{simp add: Vr-mem-f-mem, simp},$
 $\text{subst amult-pos-mono-r[of - ant (int (na n)) n-val } K \ v \ x],$
 $\text{assumption+},$
 $\text{cut-tac an-na[of } n], \text{ simp add:an-def, assumption+})$

done

lemma (**in** *Corps*) $n\text{-value-x-2-nat}:\llbracket \text{valuation } K \ v; x \in \text{carrier } (Vr \ K \ v);$
 $(an \ n) \leq ((n\text{-val } K \ v) \ x) \rrbracket \implies x \in (vp \ K \ v) \diamond (Vr \ K \ v) \ n$

by ($\text{frule n-value-x-2[of } v \ x \ an \ n], \text{ assumption+},$
 $\text{simp, simp add:r-apow-def},$
 $\text{case-tac } an \ n = \infty, \text{ simp add:an-def, simp},$
 $\text{case-tac } n = 0, \text{ simp},$
 $\text{subgoal-tac } an \ n \neq 0, \text{ simp, simp add:na-an},$
 $\text{rule contrapos-pp, simp+}, \text{ simp add:an-def})$

lemma (**in** *Corps*) $n\text{-val-n-pow}:\llbracket \text{valuation } K \ v; x \in \text{carrier } (Vr \ K \ v); 0 \leq n \rrbracket \implies$
 $(n \leq (n\text{-val } K \ v \ x)) = (x \in (vp \ K \ v) \ (Vr \ K \ v) \ n)$

by ($\text{rule iffI, simp add:n-value-x-2, simp add:n-value-x-1})$

lemma (**in** *Corps*) $\text{equal-in-vpr-apow}:\llbracket \text{valuation } K \ v; x \in \text{carrier } K; 0 \leq n;$
 $y \in \text{carrier } K; n\text{-val } K \ v \ x = n\text{-val } K \ v \ y; x \in (vp \ K \ v) \ (Vr \ K \ v) \ n \rrbracket \implies$
 $y \in (vp \ K \ v) \ (Vr \ K \ v) \ n$

apply ($\text{frule n-value-x-1[of } v \ n \ x], \text{ assumption+}, \text{ simp},$

$\text{rule } n\text{-value-}x\text{-}2[\text{of } v \ y \ n], \text{ assumption+},$
 $\text{frule } \text{mem-vp-apow-mem-Vr}[\text{of } v \ n \ x], \text{ assumption+})$
apply ($\text{frule } \text{val-pos-mem-Vr}[\text{THEN sym, of } v \ x], \text{ assumption+}, \text{ simp},$
 $\text{simp add:val-pos-n-val-pos}[\text{of } v \ x],$
 $\text{simp add:val-pos-n-val-pos}[\text{THEN sym, of } v \ y],$
 $\text{simp add:val-pos-mem-Vr, assumption+})$
done

lemma (**in** *Corps*) *convergenceTr*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; b \in \text{carrier } K;$
 $b \in (\text{vp } K \ v)(\text{Vr } K \ v) \ n; (\text{Abs } (n\text{-val } K \ v \ x)) \leq n \rrbracket \implies$
 $x \cdot_r b \in (\text{vp } K \ v)(\text{Vr } K \ v) (n + (n\text{-val } K \ v \ x))$

apply ($\text{cut-tac Abs-pos}[\text{of } n\text{-val } K \ v \ x],$
 $\text{frule ale-trans}[\text{of } 0 \ \text{Abs } (n\text{-val } K \ v \ x) \ n], \text{ assumption+},$
 $\text{thin-tac } 0 \leq \text{Abs } (n\text{-val } K \ v \ x))$
apply ($\text{frule Vr-ring}[\text{of } v],$
 $\text{frule-tac aadd-le-mono}[\text{of Abs } (n\text{-val } K \ v \ x) \ n \ n\text{-val } K \ v \ x],$
 $\text{cut-tac Abs-x-plus-x-pos}[\text{of } n\text{-val } K \ v \ x],$
 $\text{frule ale-trans}[\text{of } 0 \ \text{Abs } (n\text{-val } K \ v \ x) + n\text{-val } K \ v \ x$
 $n + n\text{-val } K \ v \ x], \text{ assumption+},$
 $\text{thin-tac } 0 \leq \text{Abs } (n\text{-val } K \ v \ x) + n\text{-val } K \ v \ x,$
 $\text{thin-tac Abs } (n\text{-val } K \ v \ x) + n\text{-val } K \ v \ x \leq n + n\text{-val } K \ v \ x,$
 $\text{rule } n\text{-value-}x\text{-}2[\text{of } v \ x \cdot_r b \ n + n\text{-val } K \ v \ x], \text{ assumption+})$
apply ($\text{frule } n\text{-value-}x\text{-}1[\text{of } v \ n \ b], \text{ assumption+})$
apply ($\text{frule aadd-le-mono}[\text{of } n \ n\text{-val } K \ v \ b \ n\text{-val } K \ v \ x],$
 $\text{frule ale-trans}[\text{of } 0 \ n + n\text{-val } K \ v \ x \ n\text{-val } K \ v \ b + n\text{-val } K \ v \ x],$
 $\text{assumption})$
apply ($\text{thin-tac } 0 \leq n + n\text{-val } K \ v \ x,$
 $\text{thin-tac } n \leq n\text{-val } K \ v \ b,$
 $\text{thin-tac } n + n\text{-val } K \ v \ x \leq n\text{-val } K \ v \ b + n\text{-val } K \ v \ x,$
 $\text{simp add:aadd-commute}[\text{of } n\text{-val } K \ v \ b \ n\text{-val } K \ v \ x])$
apply ($\text{frule } n\text{-val-valuation}[\text{of } v],$
 $\text{simp add:val-t2p}[\text{THEN sym, of } n\text{-val } K \ v \ x \ b],$
 $\text{cut-tac field-is-ring},$
 $\text{frule Ring.ring-tOp-closed}[\text{of } K \ x \ b], \text{ assumption+},$
 $\text{simp add:val-pos-n-val-pos}[\text{THEN sym, of } v \ x \cdot_r b],$
 $\text{simp add:val-pos-mem-Vr},$
 $\text{frule } n\text{-val-valuation}[\text{of } v],$
 $\text{subst val-t2p}[\text{of } n\text{-val } K \ v], \text{ assumption+},$
 $\text{frule } n\text{-value-}x\text{-}1[\text{of } v \ n \ b], \text{ assumption+},$
 $\text{simp add:aadd-commute}[\text{of } n\text{-val } K \ v \ x \ n\text{-val } K \ v \ b],$
 $\text{rule aadd-le-mono}[\text{of } n \ n\text{-val } K \ v \ b \ n\text{-val } K \ v \ x], \text{ assumption+})$
done

lemma (**in** *Corps*) *convergenceTr1*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K;$
 $b \in (\text{vp } K \ v)(\text{Vr } K \ v) (n + \text{Abs } (n\text{-val } K \ v \ x)); 0 \leq n \rrbracket \implies$
 $x \cdot_r b \in (\text{vp } K \ v)(\text{Vr } K \ v) \ n$

apply ($\text{cut-tac field-is-ring},$
 $\text{frule Vr-ring}[\text{of } v],$

frule vp-apow-ideal[of $v \ n + \text{Abs} \ (n\text{-val} \ K \ v \ x)$],
cut-tac Abs-pos[of $n\text{-val} \ K \ v \ x$],
rule aadd-two-pos[of $n \ \text{Abs} \ (n\text{-val} \ K \ v \ x)$], *assumption+*)

apply (*frule Ring.ideal-subset*[of $Vr \ K \ v \ vp \ K \ v \ (Vr \ K \ v) \ (n + \text{Abs} \ (n\text{-val} \ K \ v \ x))$], *assumption+*,
frule Vr-mem-f-mem[of $v \ b$], *assumption*,
frule convergenceTr[of $v \ x \ b \ n + \text{Abs} \ (n\text{-val} \ K \ v \ x)$], *assumption+*,
rule aadd-pos-le[of $n \ \text{Abs} \ (n\text{-val} \ K \ v \ x)$], *assumption*)

apply (*frule apos-in-aug-inf*[of n],
cut-tac Abs-pos[of $n\text{-val} \ K \ v \ x$],
frule apos-in-aug-inf[of $\text{Abs} \ (n\text{-val} \ K \ v \ x)$],
frule n-value-in-aug-inf[of $v \ x$], *assumption+*,
frule aadd-assoc-i[of $n \ \text{Abs} \ (n\text{-val} \ K \ v \ x) \ n\text{-val} \ K \ v \ x$],
assumption+,
cut-tac Abs-x-plus-x-pos[of $n\text{-val} \ K \ v \ x$])

apply (*frule-tac Ring.ring-tOp-closed*[of $K \ x \ b$], *assumption+*,
rule n-value-x-2[of $v \ x \cdot_r \ b \ n$], *assumption+*)

apply (*subst val-pos-mem-Vr*[*THEN sym*, of $v \ x \cdot_r \ b$], *assumption+*,
subst val-pos-n-val-pos[of $v \ x \cdot_r \ b$], *assumption+*)

apply (*frule n-value-x-1*[of $v \ n + \text{Abs}(n\text{-val} \ K \ v \ x) + n\text{-val} \ K \ v \ x \cdot_r \ b$],
subst aadd-assoc-i, *assumption+*,
rule aadd-two-pos[of n], *assumption+*,
rule ale-trans[of $0 \ n + \text{Abs} \ (n\text{-val} \ K \ v \ x) + n\text{-val} \ K \ v \ x$
 $n\text{-val} \ K \ v \ (x \cdot_r \ b)$],
simp, *simp add:aadd-two-pos*, *assumption*,
frule n-value-x-1[of $v \ n + \text{Abs} \ (n\text{-val} \ K \ v \ x) \ b$],
cut-tac Abs-pos[of $n\text{-val} \ K \ v \ x$],
rule aadd-two-pos[of $n \ \text{Abs} \ (n\text{-val} \ K \ v \ x)$], *assumption+*)

apply (*frule n-val-valuation*[of v],
subst val-t2p[of $n\text{-val} \ K \ v$], *assumption+*)

apply (*frule aadd-le-mono*[of $n + \text{Abs} \ (n\text{-val} \ K \ v \ x) \ n\text{-val} \ K \ v \ b$
 $n\text{-val} \ K \ v \ x$],
simp add:aadd-commute[of $n\text{-val} \ K \ v \ b \ n\text{-val} \ K \ v \ x$],
rule ale-trans[of $n \ n + (\text{Abs} \ (n\text{-val} \ K \ v \ x) + n\text{-val} \ K \ v \ x)$
 $n\text{-val} \ K \ v \ x + n\text{-val} \ K \ v \ b$],
frule aadd-pos-le[of $\text{Abs} \ (n\text{-val} \ K \ v \ x) + n\text{-val} \ K \ v \ x \ n$],
simp add:aadd-commute[of n], *assumption+*)

done

lemma (**in** *Corps*) *vp-potent-zero*: $\llbracket \text{valuation} \ K \ v; 0 \leq n \rrbracket \implies$
 $(n = \infty) = (vp \ K \ v \ (Vr \ K \ v) \ n = \{\mathbf{0}_{Vr \ K \ v}\})$

apply (*rule iffI*)

apply (*simp add:r-apow-def*, *rule contrapos-pp*, *simp+*,

frule apos-neq-minf[of n],
cut-tac mem-ant[of n], *simp*, *erule exE*, *simp*,
simp add:ant-0[*THEN sym*], *thin-tac* $n = \text{ant } z$)

apply (*case-tac* $z = 0$, *simp add:ant-0*, *simp add:r-apow-def*,
frule Vr-ring[of v],
frule Ring.ring-one[of $Vr \ K \ v$], *simp*,
simp add:Vr-0-f-0, *simp add:Vr-1-f-1*,
frule value-of-one[of v], *simp*, *simp add:value-of-zero*,
cut-tac $n = z$ **in** *zneg-aneq*[of $- \ 0$], *simp only:ant-0*)
apply (*simp add:r-apow-def*,
frule-tac $n = na \ (\text{ant } z)$ **in** *n-value-idealTr*[of v],
simp add:na-def,
simp, *thin-tac* $vp \ K \ v \ \Diamond (Vr \ K \ v) \ (na \ (\text{ant } z)) = \{0_{Vr \ K \ v}\}$,
frule Vr-ring[of v],
frule Pg-in-Vr[of v],
frule-tac $n = na \ (\text{ant } z)$ **in** *Ring.npClose*[of $Vr \ K \ v \ Pg \ K \ v$],
assumption)
apply (*frule-tac* $a = (Pg \ K \ v) \ \neg (Vr \ K \ v) \ (na \ (\text{ant } z))$ **in**
Ring.a-in-principal[of $Vr \ K \ v$], *assumption*,
simp, *frule Vr-integral*[of v],
frule val-Pg[of v], *simp*, (*erule conjE*) $+$,
frule-tac $n = na \ (\text{ant } z)$ **in** *Idomain.idom-potent-nonzero*[of $Vr \ K \ v$
Pg \ K \ v], *assumption* $+$,
simp add:Vr-0-f-0, *simp*)
done

lemma (**in** *Corps*) *Vr-potent-eqTr1*: $\llbracket \text{valuation } K \ v; \ 0 \leq n; \ 0 \leq m;$
 $(vp \ K \ v) \ (Vr \ K \ v) \ n = (vp \ K \ v) \ (Vr \ K \ v) \ m; \ m = 0 \rrbracket \implies n = m$

apply (*frule Vr-ring*[of v],
simp add:r-apow-def,
case-tac $n = 0$, *simp*,
case-tac $n = \infty$, *simp*,
frule val-Pg[of v], *erule conjE*, *simp*,
erule conjE,
rotate-tac -3 , *drule sym*,
frule Lv-pos[of v], *simp*,
frule val-poss-mem-Vr[of $v \ Pg \ K \ v$], *assumption* $+$,
drule sym, *simp*, *simp add:Vr-0-f-0*)

apply (*simp*,
drule sym,
frule Ring.ring-one[of $Vr \ K \ v$], *simp*,

frule n-value-x-1-nat[of $v \ 1_r (Vr \ K \ v) \ na \ n$], *assumption*,
simp add:an-na, *simp add:Vr-1-f-1*,
frule n-val-valuation[of v],

$\text{simp add: value-of-one[of } n\text{-val } K \text{ } v])$
done

lemma (in Corps) Vr-potent-eqTr2: $\llbracket \text{valuation } K \text{ } v; (vp \text{ } K \text{ } v) \Diamond (Vr \text{ } K \text{ } v) \text{ } n = (vp \text{ } K \text{ } v) \Diamond (Vr \text{ } K \text{ } v) \text{ } m \rrbracket \implies n = m$

apply (frule Vr-ring[of v],
frule val-Pg[of v], simp, (erule conjE)+,
rotate-tac -1, frule sym, thin-tac v (Pg K v) = Lv K v,
frule Lv-pos[of v], simp)

apply (subgoal-tac 0 ≤ int n, subgoal-tac 0 ≤ int m,
frule n-value-idealTr[of v m]) **apply** simp **apply** simp
apply (
thin-tac vp K v $\Diamond (Vr \text{ } K \text{ } v) \text{ } m = Vr \text{ } K \text{ } v \Diamond_p (Pg \text{ } K \text{ } v \neg (Vr \text{ } K \text{ } v) \text{ } m)$,
frule n-value-idealTr[of v n], simp, simp,
thin-tac vp K v $\Diamond (Vr \text{ } K \text{ } v) \text{ } n = Vr \text{ } K \text{ } v \Diamond_p (Pg \text{ } K \text{ } v \neg (Vr \text{ } K \text{ } v) \text{ } m)$,
frule val-poss-mem-Vr[of v Pg K v], assumption+)

apply (frule Lv-z[of v], erule exE,
rotate-tac -4, drule sym, simp,
frule eq-ideal-eq-val[of v Pg K v $\neg (Vr \text{ } K \text{ } v) \text{ } n$ Pg K v $\neg (Vr \text{ } K \text{ } v) \text{ } m$])
apply (rule Ring.npClose, assumption+, rule Ring.npClose, assumption+)
apply (simp only: Vr-exp-f-exp,
simp add: val-exp-ring[THEN sym, of v Pg K v],
thin-tac Vr K v $\Diamond_p (Pg \text{ } K \text{ } v \neg^K n) = Vr \text{ } K \text{ } v \Diamond_p (Pg \text{ } K \text{ } v \neg^K m)$)

apply (case-tac n = 0, simp, case-tac m = 0, simp,
simp only: of-nat-0-less-iff[THEN sym, of m],
simp only: asprod-amult a-z-z,
simp only: ant-0[THEN sym], simp only: aeq-zeq, simp)
apply (auto simp add: asprod-mult)
done

lemma (in Corps) Vr-potent-eq: $\llbracket \text{valuation } K \text{ } v; 0 \leq n; 0 \leq m; (vp \text{ } K \text{ } v) (Vr \text{ } K \text{ } v) \text{ } n = (vp \text{ } K \text{ } v) (Vr \text{ } K \text{ } v) \text{ } m \rrbracket \implies n = m$

apply (frule n-val-valuation[of v],
case-tac m = 0,
simp add: Vr-potent-eqTr1)
apply (case-tac n = 0,
frule sym, thin-tac vp K v (Vr K v) n = vp K v (Vr K v) m,
frule Vr-potent-eqTr1[of v m n], assumption+,
rule sym, assumption,
frule vp-potent-zero[of v n], assumption+)

apply (case-tac n = ∞, simp,
thin-tac vp K v (Vr K v) ∞ = {0}_{Vr K v},

frule vp-potent-zero[THEN sym, of v m], assumption+, simp,
simp,
frule vp-potent-zero[THEN sym, of v m], assumption+, simp,
thin-tac vp K v (Vr K v) m ≠ {0_{Vr K v}})

apply (*frule aposs-na-poss[of n], subst less-le, simp,*
frule aposs-na-poss[of m], subst less-le, simp,
simp add:r-apow-def,
frule Vr-potent-eqTr2[of v na n na m], assumption+,
thin-tac vp K v ◇(Vr K v) (na n) = vp K v ◇(Vr K v) (na m),
simp add:aeq-nat-eq[THEN sym])

done

the following two lemma (in Corps) s are used in completion of K

lemma (in Corps) *Vr-prime-maximalTr1: [valuation K v; x ∈ carrier (Vr K v);*
Suc 0 < n] ⇒ x ·_r(Vr K v) (x^{~K} (n - Suc 0)) ∈ (Vr K v) ◇_p (x^{~K} n)

apply (*frule Vr-ring[of v],*
subgoal-tac x^{~K} n = x^{~K} (Suc (n - Suc 0)),
simp del:Suc-pred,
rotate-tac -1, drule sym)

apply (*subst Vr-tOp-f-tOp, assumption+,*
subst Vr-exp-f-exp[of v, THEN sym], assumption+,
simp only:Ring.npClose, simp del:Suc-pred)

apply (*cut-tac field-is-ring,*
frule Ring.npClose[of K x n - Suc 0],
frule Vr-mem-f-mem[of v x], assumption+,
frule Vr-mem-f-mem[of v x], assumption+)
apply (*simp add:Ring.ring-tOp-commute[of K x x^{~K} (n - Suc 0)]*)

apply (*rule Ring.a-in-principal, assumption*)

apply (*frule Ring.npClose[of Vr K v x n], assumption,*
simp add:Vr-exp-f-exp)

apply (*simp only:Suc-pred*)

done

lemma (in Corps) *Vr-prime-maximalTr2: [valuation K v; x ∈ vp K v; x ≠ 0;*
Suc 0 < n] ⇒ x ∉ Vr K v ◇_p (x^{~K} n) ∧ x^{~K} (n - Suc 0) ∉ (Vr K v) ◇_p
(x^{~K} n)

apply (*frule Vr-ring[of v]*)

apply (*frule vp-mem-Vr-mem[of v x], assumption,*
frule Ring.npClose[of Vr K v x n],
simp only:Vr-exp-f-exp)

apply (*cut-tac field-is-ring,*
cut-tac field-is-idom,
frule Vr-mem-f-mem[of v x], assumption+,
frule Idomain.idom-potent-nonzero[of K x n], assumption+)

apply (*rule conjI*)

apply (*rule contrapos-pp, simp+*)

apply (*frule val-Rxa-gt-a[of v x^{~K} n x],*

```

      simp, simp add: Vr-exp-f-exp, assumption+)
apply (simp add: val-exp-ring[THEN sym, of v x n])
apply (frule val-nonzero-z[of v x], assumption+, erule exE,
      simp add: asprod-amult a-z-z)
apply (simp add: vp-mem-val-poss[of v x])
apply (rule contrapos-pp, simp+)
apply (frule val-Rxa-gt-a[of v xK n xK (n - Suc 0)])
      apply (simp, frule Ring.npClose[of Vr K v x n - Suc 0], assumption+)
      apply (simp add: Vr-exp-f-exp)
      apply (frule Ring.npClose[of Vr K v x n - Suc 0], assumption+,
      simp add: Vr-exp-f-exp, assumption)
apply (simp add: val-exp-ring[THEN sym, of v x])
apply (simp add: vp-mem-val-poss[of v x])
apply (frule val-nonzero-z[of v x], assumption+, erule exE,
      simp add: asprod-amult a-z-z)
done

lemma (in Corps) Vring-prime-maximal: [[valuation K v; prime-ideal (Vr K v) I;
  I ≠ {0Vr K v}]] ⇒ maximal-ideal (Vr K v) I
apply (frule Vr-ring[of v],
  frule Ring.prime-ideal-proper[of Vr K v I], assumption+,
  frule Ring.prime-ideal-ideal[of Vr K v I], assumption+,
  frule ideal-pow-vp[of v I],
  frule n-value-idealTr[of v na (n-val K v (Ig K v I))],
  simp, simp, assumption+)

apply (case-tac na (n-val K v (Ig K v I)) = 0,
  simp, frule Ring.Rxa-one[of Vr K v], simp,
  frule Suc-leI[of 0 na (n-val K v (Ig K v I))],
  thin-tac 0 < na (n-val K v (Ig K v I)))
apply (case-tac na (n-val K v (Ig K v I)) = Suc 0, simp,
  frule Pg-in-Vr[of v])
apply (frule vp-maximal[of v],
  frule Ring.maximal-ideal-ideal[of Vr K v vp K v], assumption+,
  subst Ring.idealprod-whole-r[of Vr K v vp K v], assumption+)

apply (rotate-tac -1, drule not-sym,
  frule le-neq-implies-less[of Suc 0 na (n-val K v (Ig K v I))],
  assumption+,
  thin-tac Suc 0 ≤ na (n-val K v (Ig K v I)),
  thin-tac Suc 0 ≠ na (n-val K v (Ig K v I)),
  thin-tac Vr K v ⋄p 1r Vr K v = carrier (Vr K v))
apply (frule val-Pg[of v], simp, (erule conjE)+,
  frule Lv-pos[of v], rotate-tac -2, drule sym)
apply (frule val-poss-mem-Vr[of v Pg K v],
  frule vp-mem-val-poss[THEN sym, of v Pg K v], assumption+, simp)

apply (frule Vr-prime-maximalTr2[of v Pg K v
  na (n-val K v (Ig K v I))],

```

```

      simp add:vp-mem-val-poss[of v Pg K v], assumption+, erule conjE)
apply (frule Ring.npMulDistr[of Vr K v Pg K v na 1 na (n-val K v (Ig K v I)) -
Suc 0], assumption+, simp add:na-1)

apply (rotate-tac 8, drule sym)
apply (frule Ring.a-in-principal[of Vr K v
Pg K v  $\neg$ (Vr K v) (na (n-val K v (Ig K v I)))], simp add:Ring.npClose)

apply (simp add:Vr-exp-f-exp[of v])
  apply (simp add:Ring.ring-l-one[of Vr K v])
  apply (frule n-value-idealTr[THEN sym,
of v na (n-val K v (Ig K v I))], simp)
  apply (simp add:Vr-exp-f-exp)
  apply (rotate-tac 6, drule sym, simp)
apply (thin-tac I  $\neq$  carrier (Vr K v),
thin-tac I = vp K v  $\Diamond$  (Vr K v) (na (n-val K v (Ig K v I))),
thin-tac v (Pg K v) = Lv K v,
thin-tac (Vr K v)  $\Diamond_p$  ((Pg K v)  $\cdot_r$  (Vr K v)
((Pg K v)  $\neg^K$  (na ((n-val K v) (Ig K v I)) - (Suc 0))) =
I,
thin-tac Pg K v  $\in$  carrier K,
thin-tac Pg K v  $\neq$  0,
thin-tac Pg K v  $\neg^K$  (na ((n-val K v) (Ig K v I))) =
Pg K v  $\cdot_r$  Vr K v Pg K v  $\neg^K$  ((na ((n-val K v) (Ig K v I))) - Suc 0))

apply (simp add:prime-ideal-def, erule conjE,
drule-tac x = Pg K v in bspec, assumption,
drule-tac x = Pg K v  $\neg^K$  (na (n-val K v (Ig K v I)) - Suc 0) in bspec)
  apply (simp add:Vr-exp-f-exp[THEN sym, of v])
apply (rule Ring.npClose[of Vr K v Pg K v], assumption+)
apply simp
done

```

From the above lemma (in Corps) , we see that a valuation ring is of dimension one.

```

lemma (in Corps) field-frac1:  $\llbracket 1_r \neq 0; x \in \text{carrier } K \rrbracket \implies x = x \cdot_r ((1_r)^{-K})$ 
by (simp add:invf-one,
cut-tac field-is-ring,
simp add:Ring.ring-r-one[THEN sym])

```

```

lemma (in Corps) field-frac2:  $\llbracket x \in \text{carrier } K; x \neq 0 \rrbracket \implies x = (1_r) \cdot_r ((x^{-K})^{-K})$ 
by (cut-tac field-is-ring, simp add:field-inv-inv,
simp add:Ring.ring-l-one[THEN sym])

```

```

lemma (in Corps) val-nonpos-inv-pos:  $\llbracket \text{valuation } K \text{ v}; x \in \text{carrier } K; \neg 0 \leq (v \ x) \rrbracket \implies 0 < (v \ (x^{-K}))$ 
by (case-tac x = 0_K, simp add:value-of-zero,

```

```

    frule Vr-ring[of v],
    simp add:aneg-le[of 0 v x],
    frule value-of-inv[THEN sym, of v x], assumption+,
    frule aless-minus[of v x 0], simp)

lemma (in Corps) frac-Vr-is-K: [valuation K v; x ∈ carrier K] ⇒
  ∃ s ∈ carrier (Vr K v). ∃ t ∈ carrier (Vr K v) - {0}. x = s ·r (t-K)
apply (frule Vr-ring[of v],
      frule has-val-one-neq-zero[of v])
apply (case-tac x = 0K,
      frule Ring.ring-one[of Vr K v],
      frule field-frac1[of x],
      simp only: Vr-1-f-1, frule Ring.ring-zero[of Vr K v],
      simp add: Vr-0-f-0 Vr-1-f-1, blast)
apply (case-tac 0 ≤ (v x),
      frule val-pos-mem-Vr[THEN sym, of v x], assumption+, simp,
      frule field-frac1[of x], assumption+,
      frule has-val-one-neq-zero[of v],
      frule Ring.ring-one[of Vr K v], simp only: Vr-1-f-1, blast)
apply (frule val-nonpos-inv-pos[of v x], assumption+,
      cut-tac invf-inv[of x], erule conjE,
      frule val-poss-mem-Vr[of v x-K], assumption+)
apply (frule Ring.ring-one[of Vr K v], simp only: Vr-1-f-1,
      frule field-frac2[of x], assumption+)
apply (cut-tac invf-closed1[of x], blast, simp+)
done

lemma (in Corps) valuations-eqTr1: [valuation K v; valuation K v';
  Vr K v = Vr K v'; ∀ x ∈ carrier (Vr K v). v x = v' x] ⇒ v = v'
apply (rule funcset-eq [of - carrier K],
      simp add: valuation-def, simp add: valuation-def,
      rule ballI,
      frule-tac x = x in frac-Vr-is-K[of v], assumption+,
      (erule bexE)+, simp, erule conjE)
apply (frule-tac x = t in Vr-mem-f-mem[of v'], assumption,
      cut-tac x = t in invf-closed1, simp, simp, erule conjE)
  apply (frule-tac x = s in Vr-mem-f-mem[of v'], assumption+,
        simp add: val-t2p, simp add: value-of-inv)
done

lemma (in Corps) ridmap-rhom: [ valuation K v; valuation K v';
  carrier (Vr K v) ⊆ carrier (Vr K v') ] ⇒
  ridmap (Vr K v) ∈ rHom (Vr K v) (Vr K v')
apply (frule Vr-ring[of v], frule Vr-ring[of v'],
      subst rHom-def, simp, rule conjI)
apply (simp add: aHom-def, rule conjI,
      rule Pi-I, simp add: ridmap-def subsetD,
      simp add: ridmap-def restrict-def extensional-def,
      (rule ballI)+,

```


$\text{frule Ring.ring-is-ag[of Vr K v], simp add:aGroup.ag-pOp-closed,}$
 $\text{simp add:Vr-pOp-f-pOp subsetD)}$
apply (rule conjI, (rule ballI)+, simp add:ridmap-def,
 $\text{simp add:Ring.ring-tOp-closed, simp add:Vr-tOp-f-tOp subsetD,}$
 $\text{frule Ring.ring-one[of Vr K v], frule Ring.ring-one[of Vr K v'],}$
 $\text{simp add:Vr-1-f-1, simp add:ridmap-def)}$
done

lemma (in Corps) contract-ideal: $\llbracket \text{valuation K v; valuation K v';}$
 $\text{carrier (Vr K v) } \subseteq \text{carrier (Vr K v')} \rrbracket \implies$
 $\text{ideal (Vr K v) (carrier (Vr K v) } \cap \text{vp K v')}$
apply (frule-tac ridmap-rhom[of v v'], assumption+,
 $\text{frule Vr-ring[of v], frule Vr-ring[of v']})$
apply (cut-tac TwoRings.i-contract-ideal[of Vr K v Vr K v'
 $\text{ridmap (Vr K v) vp K v'},$
 $\text{subgoal-tac (i-contract (ridmap (Vr K v)) (Vr K v) (Vr K v')}$
 $\text{(vp K v')) = (carrier (Vr K v) } \cap \text{vp K v'))}$
apply simp
apply (thin-tac ideal (Vr K v) (i-contract (ridmap (Vr K v))
 $\text{(Vr K v) (Vr K v') (vp K v')},$
 $\text{simp add:i-contract-def invim-def ridmap-def, blast})$
apply (simp add:TwoRings-def TwoRings-axioms-def, simp)
apply (simp add:vp-ideal)
done

lemma (in Corps) contract-prime: $\llbracket \text{valuation K v; valuation K v';}$
 $\text{carrier (Vr K v) } \subseteq \text{carrier (Vr K v')} \rrbracket \implies$
 $\text{prime-ideal (Vr K v) (carrier (Vr K v) } \cap \text{vp K v')}$
apply (frule-tac ridmap-rhom[of v v'], assumption+,
 $\text{frule Vr-ring[of v],}$
 $\text{frule Vr-ring[of v'],}$
 $\text{cut-tac TwoRings.i-contract-prime[of Vr K v Vr K v' ridmap (Vr K v)}$
 vp K v')
apply (subgoal-tac (i-contract (ridmap (Vr K v)) (Vr K v) (Vr K v')
 $\text{(vp K v')) = (carrier (Vr K v) } \cap \text{vp K v')},$
 simp,
 $\text{thin-tac prime-ideal (Vr K v) (i-contract}$
 $\text{(ridmap (Vr K v)) (Vr K v) (Vr K v') (vp K v')},$
 $\text{simp add:i-contract-def invim-def ridmap-def, blast})$
apply (simp add:TwoRings-def TwoRings-axioms-def, simp)
apply (simp add:vp-prime)
done

lemma (in Corps) valuation-equivTr: $\llbracket \text{valuation K v; valuation K v';}$
 $\text{x } \in \text{carrier K; } 0 < (v' x); \text{carrier (Vr K v) } \subseteq \text{carrier (Vr K v')} \rrbracket$
 $\implies 0 \leq (v x)$
apply (rule contrapos-pp, simp+,
 $\text{frule val-nonpos-inv-pos[of v x], assumption+},$

$\text{case-tac } x = \mathbf{0}_K, \text{ simp add:value-of-zero[of } v] \text{) apply (}$
 $\text{cut-tac invf-closed1[of } x], \text{ simp, erule conjE,}$
 $\text{frule aless-imp-le[of } 0 \text{ } v \text{ } (x^{-K})])$
apply ($\text{simp add:val-pos-mem-Vr[of } v \text{ } x^{-K}],$
 $\text{frule subsetD[of carrier (Vr } K \text{ } v) \text{ carrier (Vr } K \text{ } v') \text{ } x^{-K}],$
 assumption+,
 $\text{frule val-pos-mem-Vr[THEN sym, of } v' \text{ } x^{-K}], \text{ assumption+})$
apply ($\text{simp, simp add:value-of-inv[of } v' \text{ } x],$
 $\text{cut-tac ale-minus[of } 0 - v' \text{ } x], \text{ thin-tac } 0 \leq -v' \text{ } x,$
 $\text{simp only:a-minus-minus,}$
 $\text{cut-tac aneg-less[THEN sym, of } v' \text{ } x - 0], \text{ simp,}$
 $\text{assumption, simp})$
done

lemma (**in** *Corps*) *contract-maximal*: $[\text{valuation } K \text{ } v; \text{ valuation } K \text{ } v';$
 $\text{carrier (Vr } K \text{ } v) \subseteq \text{carrier (Vr } K \text{ } v') \implies$
 $\text{maximal-ideal (Vr } K \text{ } v) (\text{carrier (Vr } K \text{ } v) \cap \text{vp } K \text{ } v')]$
apply ($\text{frule Vr-ring[of } v],$
 $\text{frule Vr-ring[of } v'],$
 $\text{rule Vring-prime-maximal, assumption+,}$
 $\text{simp add:contract-prime})$
apply ($\text{frule vp-nonzero[of } v'],$
 $\text{frule vp-ideal[of } v'],$
 $\text{frule Ring.ideal-zero[of Vr } K \text{ } v' \text{ vp } K \text{ } v'], \text{ assumption+,}$
 $\text{frule sets-not-eq[of vp } K \text{ } v' \{ \mathbf{0}_{(\text{Vr } K \text{ } v')} \}],$
 $\text{simp add: singleton-sub[of } \mathbf{0}_{(\text{Vr } K \text{ } v')} \text{ carrier (Vr } K \text{ } v')],$
 $\text{erule bexE, simp add:Vr-0-f-0})$

apply ($\text{case-tac } a \in \text{carrier (Vr } K \text{ } v), \text{ blast,}$
 $\text{frule-tac } x = a \text{ in vp-mem-Vr-mem[of } v'], \text{ assumption+,}$
 $\text{frule-tac } x = a \text{ in Vr-mem-f-mem[of } v'], \text{ assumption+,}$
 $\text{subgoal-tac } a \in \text{carrier (Vr } K \text{ } v), \text{ blast,}$
 $\text{frule-tac } x1 = a \text{ in val-pos-mem-Vr[THEN sym, of } v], \text{ assumption+,}$
 $\text{simp, frule val-nonpos-inv-pos[of } v], \text{ assumption+})$

apply ($\text{frule-tac } y = v \text{ } (a^{-K}) \text{ in aless-imp-le[of } 0],$
 $\text{cut-tac } x = a \text{ in invf-closed1, simp,}$
 $\text{frule-tac } x = a^{-K} \text{ in val-poss-mem-Vr[of } v], \text{ simp, assumption+})$
apply ($\text{frule-tac } c = a^{-K} \text{ in subsetD[of carrier (Vr } K \text{ } v)$
 $\text{carrier (Vr } K \text{ } v')], \text{ assumption+}) \text{ apply (}$
 $\text{frule-tac } x = a^{-K} \text{ in val-pos-mem-Vr[of } v'],$
 $\text{simp, simp only:value-of-inv[of } v'], \text{ simp,}$
 $\text{simp add:value-of-inv[of } v'])$

apply ($\text{frule-tac } y = -v' \text{ } a \text{ in ale-minus[of } 0], \text{ simp add:a-minus-minus,}$
 $\text{frule-tac } x = a \text{ in vp-mem-val-poss[of } v'], \text{ assumption+,}$
 $\text{simp})$
done

2.6 Equivalent valuations

definition

$v\text{-equiv} :: [-, 'r \Rightarrow \text{ant}, 'r \Rightarrow \text{ant}] \Rightarrow \text{bool}$ **where**
 $v\text{-equiv } K \ v1 \ v2 \longleftrightarrow n\text{-val } K \ v1 = n\text{-val } K \ v2$

lemma (in Corps) valuation-equivTr1: $[[\text{valuation } K \ v; \text{valuation } K \ v';$
 $\forall x \in \text{carrier } K. 0 \leq (v \ x) \longrightarrow 0 \leq (v' \ x)]] \implies$
 $\text{carrier } (Vr \ K \ v) \subseteq \text{carrier } (Vr \ K \ v')$

apply (frule Vr-ring[of v],
frule Vr-ring[of v'])
apply (rule subsetI,
case-tac $x = \mathbf{0}_K$, simp, simp add: Vr-def Sr-def,
frule-tac $x1 = x$ in val-pos-mem-Vr[THEN sym, of v],
frule-tac $x = x$ in Vr-mem-f-mem[of v],
simp, frule-tac $x = x$ in Vr-mem-f-mem[of v], assumption+)
apply (drule-tac $x = x$ in bspec, simp add: Vr-mem-f-mem)
apply simp
apply (subst val-pos-mem-Vr[THEN sym, of v'], assumption+,
simp add: Vr-mem-f-mem, assumption+)
done

lemma (in Corps) valuation-equivTr2: $[[\text{valuation } K \ v; \text{valuation } K \ v';$
 $\text{carrier } (Vr \ K \ v) \subseteq \text{carrier } (Vr \ K \ v'); \text{vp } K \ v = \text{carrier } (Vr \ K \ v) \cap \text{vp } K \ v]]$
 $\implies \text{carrier } (Vr \ K \ v') \subseteq \text{carrier } (Vr \ K \ v)$

apply (frule Vr-ring[of v], frule Vr-ring[of v'])
apply (rule subsetI)
apply (case-tac $x = \mathbf{0}_{(Vr \ K \ v')}$, simp,
subst Vr-0-f-0[of v'], assumption+,
subst Vr-0-f-0[of v, THEN sym], assumption,
simp add: Ring.ring-zero)
apply (rule contrapos-pp, simp+)
apply (frule-tac $x = x$ in Vr-mem-f-mem[of v'], assumption+)
apply (simp add: val-pos-mem-Vr[THEN sym, of v])
apply (cut-tac $y = v \ x$ in aneg-le[of 0], simp)
apply (simp add: Vr-0-f-0[of v'])
apply (frule-tac $x = v \ x$ in aless-minus[of - 0], simp,
thin-tac $v \ x < 0$, thin-tac $\neg 0 \leq v \ x$)
apply (simp add: value-of-inv[THEN sym, of v])
apply (cut-tac $x = x$ in invf-closed1, simp, simp, erule conjE)
apply (frule-tac $x1 = x^{-K}$ in vp-mem-val-poss[THEN sym, of v],
assumption, simp, erule conjE)
apply (frule vp-ideal [of v'])
apply (frule-tac $x = x^{-K}$ and $r = x$ in Ring.ideal-ring-multiple[of Vr K v'
vp K v'], assumption+)
apply (frule-tac $x = x^{-K}$ in vp-mem-Vr-mem[of v'], assumption+)
apply (frule-tac $x = x$ and $y = x^{-K}$ in Ring.ring-tOp-commute[of Vr K v'],
assumption+, simp,

```

      thin-tac x ·r Vr K v' x- K = x- K ·r Vr K v' x)
apply (simp add: Vr-tOp-f-tOp)
apply (cut-tac x = x in linvf, simp, simp)
apply (cut-tac field-is-ring, frule Ring.ring-one[of K])
apply (frule ideal-inc-elem0val-whole[of v' 1r vp K v'],
      assumption+, simp add:value-of-one, assumption+)
apply (frule vp-not-whole[of v'], simp)
done

lemma (in Corps) eq-carr-eq-Vring:  $\llbracket \text{valuation } K \ v; \text{valuation } K \ v';$ 
       $\text{carrier } (Vr \ K \ v) = \text{carrier } (Vr \ K \ v') \rrbracket \implies Vr \ K \ v = Vr \ K \ v'$ 
apply (simp add: Vr-def Sr-def)
done

lemma (in Corps) valuations-equiv:  $\llbracket \text{valuation } K \ v; \text{valuation } K \ v';$ 
       $\forall x \in \text{carrier } K. 0 \leq (v \ x) \longrightarrow 0 \leq (v' \ x) \rrbracket \implies v\text{-equiv } K \ v \ v'$ 

apply (frule Vr-ring[of v], frule Vr-ring[of v'])

apply (frule valuation-equivTr1[of v v'], assumption+)

apply (frule contract-maximal [of v v'], assumption+)

apply (frule Vr-local[of v (carrier (Vr K v)  $\cap$  vp K v')],
      assumption+)

apply (frule valuation-equivTr2[of v v'], assumption+,
      frule equalityI[of carrier (Vr K v) carrier (Vr K v')],
      assumption+,
      thin-tac carrier (Vr K v)  $\subseteq$  carrier (Vr K v'),
      thin-tac carrier (Vr K v')  $\subseteq$  carrier (Vr K v))

apply (frule vp-ideal[of v'],
      frule Ring.ideal-subset1[of Vr K v' vp K v'], assumption,
      simp add: Int-absorb1,
      thin-tac  $\forall x \in \text{carrier } K. 0 \leq v \ x \longrightarrow 0 \leq v' \ x$ ,
      thin-tac  $vp \ K \ v' \subseteq \text{carrier } (Vr \ K \ v')$ ,
      thin-tac ideal (Vr K v') (vp K v'),
      thin-tac maximal-ideal (Vr K v) (vp K v'))

apply (simp add: v-equiv-def,
      rule valuations-eqTr1[of n-val K v n-val K v'],
      (simp add: n-val-valuation)+,
      rule eq-carr-eq-Vring[of n-val K v n-val K v'],
      (simp add: n-val-valuation)+,

```

```

      subst Vr-n-val-Vr[THEN sym, of v], assumption+,
      subst Vr-n-val-Vr[THEN sym, of v'], assumption+)
apply (rule ballI,
  frule n-val-valuation[of v],
  frule n-val-valuation[of v'],
  frule-tac x1 = x in val-pos-mem-Vr[THEN sym, of n-val K v],
  simp add: Vr-mem-f-mem, simp,
  frule Vr-n-val-Vr[THEN sym, of v], simp,
  thin-tac carrier (Vr K (n-val K v)) = carrier (Vr K v'),
  frule-tac x1 = x in val-pos-mem-Vr[THEN sym, of v'],
    simp add: Vr-mem-f-mem,
  simp,
  frule-tac x = x in val-pos-n-val-pos[of v'],
  simp add: Vr-mem-f-mem, simp,
  frule-tac x = x in ideal-apow-n-val[of v],
  simp add: Vr-n-val-Vr[THEN sym, of v], simp)
apply (frule eq-carr-eq-Vring[of v v'], assumption+,
  frule-tac x = x in ideal-apow-n-val[of v'], assumption,
  simp add: Vr-n-val-Vr[THEN sym, of v],
  thin-tac Vr K v'  $\triangleleft_p$  x = vp K v' (Vr K v') (n-val K v x),
  frule-tac n = n-val K v' x and m = n-val K v x in
    Vr-potent-eq[of v'], assumption+,
  frule sym, assumption+)
done

lemma (in Corps) val-equiv-axiom1: valuation K v  $\implies$  v-equiv K v v
apply (simp add: v-equiv-def)
done

lemma (in Corps) val-equiv-axiom2:  $\llbracket$  valuation K v; valuation K v';
  v-equiv K v v'  $\rrbracket \implies$  v-equiv K v' v
apply (simp add: v-equiv-def)
done

lemma (in Corps) val-equiv-axiom3:  $\llbracket$  valuation K v; valuation K v';
  valuation K v'; v-equiv K v v'; v-equiv K v' v''  $\rrbracket \implies$  v-equiv K v v''
apply (simp add: v-equiv-def)
done

lemma (in Corps) n-val-equiv-val:  $\llbracket$  valuation K v  $\rrbracket \implies$ 
  v-equiv K v (n-val K v)
apply (frule valuations-equiv[of v n-val K v], simp add: n-val-valuation)
apply (rule ballI, rule impI, simp add: val-pos-n-val-pos,
  assumption)
done

```

2.7 Prime divisors

definition

$prime-divisor :: [-, 'b \Rightarrow ant] \Rightarrow$
 $('b \Rightarrow ant) \text{ set } (\langle 2P \text{ } - \rangle \rangle [96,97]96) \text{ where}$
 $P_K v = \{v'. \text{ valuation } K v' \wedge v\text{-equiv } K v v'\}$

definition

$prime-divisors :: - \Rightarrow ('b \Rightarrow ant) \text{ set set } (\langle Pds \rangle \rangle 96) \text{ where}$
 $Pds_K = \{P. \exists v. \text{ valuation } K v \wedge P = P_K v\}$

definition

$normal\text{-valuation}\text{-belonging}\text{-to}\text{-prime}\text{-divisor} ::$
 $[-, ('b \Rightarrow ant) \text{ set}] \Rightarrow ('b \Rightarrow ant) (\langle \nu - \rangle \rangle [96,97]96) \text{ where}$
 $\nu_K P = n\text{-val } K (SOME v. v \in P)$

lemma (in Corps) $val\text{-in}\text{-}P\text{-valuation}::\llbracket \text{valuation } K v; v' \in P_K v \rrbracket \Rightarrow$
 $\text{valuation } K v'$

apply (simp add:prime-divisor-def)
done

lemma (in Corps) $vals\text{-in}\text{-}P\text{-equiv}::\llbracket \text{valuation } K v; v' \in P_K v \rrbracket \Rightarrow$
 $v\text{-equiv } K v v'$

apply (simp add:prime-divisor-def)
done

lemma (in Corps) $v\text{-in}\text{-prime}\text{-}v\text{:valuation } K v \Rightarrow v \in P_K v$

apply (simp add:prime-divisor-def,
 $frule \text{ val-equiv-axiom1 [of } v], \text{ assumption+})$
done

lemma (in Corps) $some\text{-in}\text{-prime}\text{-divisor}::\text{valuation } K v \Rightarrow$
 $(SOME w. w \in P_K v) \in P_K v$

apply (subgoal-tac $P_K v \neq \{\}$,
 $rule \text{ nonempty-some[of } P_K v], \text{ assumption+},$
 $frule v\text{-in}\text{-prime}\text{-}v\text{[of } v])$

apply blast
done

lemma (in Corps) $\text{valuation}\text{-some}\text{-in}\text{-prime}\text{-divisor}::\text{valuation } K v$
 $\Rightarrow \text{valuation } K (SOME w. w \in P_K v)$

apply (frule some-in-prime-divisor[of v],
 $\text{simp add:prime-divisor-def})$

done

lemma (in Corps) $\text{valuation}\text{-some}\text{-in}\text{-prime}\text{-divisor1}::P \in Pds \Rightarrow$
 $\text{valuation } K (SOME w. w \in P)$

apply (simp add:prime-divisors-def, erule exE)
apply (simp add:valuation-some-in-prime-divisor)
done

lemma (in Corps) $\text{representative}\text{-of}\text{-pd}\text{-valuation}::$

$P \in Pds \implies \text{valuation } K (\nu_K P)$
apply (*simp add:prime-divisors-def*,
erule exE, *erule conjE*,
simp add:normal-valuation-belonging-to-prime-divisor-def,
frule-tac v = v in valuation-some-in-prime-divisor)

apply (*rule n-val-valuation*, *assumption+*)
done

lemma (*in Corps*) *some-in-P-equiv:valuation K v $\implies$*
v-equiv K v (SOME w. w $\in P_K v$)
apply (*frule some-in-prime-divisor[of v]*)
apply (*rule vals-in-P-equiv*, *assumption+*)
done

lemma (*in Corps*) *n-val-n-val1:P $\in Pds \implies$ n-val K ($\nu_K P$) = ($\nu_K P$)*
apply (*simp add: normal-valuation-belonging-to-prime-divisor-def*,
frule valuation-some-in-prime-divisor1[of P])
apply (*rule n-val-n-val[of SOME v. v $\in P$]*, *assumption+*)
done

lemma (*in Corps*) *P-eq-val-equiv:[[valuation K v; valuation K v'] $\implies$*
(v-equiv K v v') = (P_K v = P_K v')
apply (*rule iffI*,
rule equalityI,
rule subsetI, *simp add:prime-divisor-def*, *erule conjE*,
frule val-equiv-axiom2[of v v'], *assumption+*,
rule val-equiv-axiom3[of v' v], *assumption+*,
rule subsetI, *simp add:prime-divisor-def*, *erule conjE*)
apply (*rule val-equiv-axiom3[of v v']*, *assumption+*,
frule v-in-prime-v[of v], *simp*,
thin-tac P_K v = P_K v',
simp add:prime-divisor-def,
rule val-equiv-axiom2[of v' v], *assumption+*)
done

lemma (*in Corps*) *unique-n-valuation:[[P $\in Pds_K$; P' $\in Pds$] $\implies$*
(P = P') = ($\nu_K P = \nu_K P'$)
apply (*rule iffI*, *simp*)
apply (*simp add:prime-divisors-def*,
(erule exE)+, *(erule conjE)+*)
apply (*simp add:normal-valuation-belonging-to-prime-divisor-def*,
frule-tac v = v in some-in-P-equiv,
frule-tac v = va in some-in-P-equiv,
subgoal-tac v-equiv K (SOME w. w $\in P_K v$) (SOME w. w $\in P_K va$))
apply (*frule-tac v = v in some-in-prime-divisor*,
frule-tac v = va in some-in-prime-divisor,
frule-tac v = v and v' = SOME w. w $\in P_K v$ and v'' =
SOME w. w $\in P_K va$ in val-equiv-axiom3)

```

apply (simp add:prime-divisor-def,
      simp add:prime-divisor-def, assumption+,
      frule-tac  $v = va$  and  $v' = \text{SOME } w. w \in P_K va$  in
        val-equiv-axiom2,
      simp add:prime-divisor-def, assumption+)
apply (frule-tac  $v = v$  and  $v' = \text{SOME } w. w \in P_K va$  and  $v'' = va$  in
      val-equiv-axiom3,
      simp add:prime-divisor-def,
      simp add:prime-divisor-def, assumption+,
      frule-tac  $v = v$  and  $v' = va$  in  $P\text{-eq-}\text{val-equiv}$ , assumption+)
apply simp
apply (simp add:v-equiv-def)
done

```

```

lemma (in Corps)  $n\text{-val-representative}: P \in Pds \implies (\nu_K P) \in P$ 
apply (simp add:prime-divisors-def,
      erule exE, erule conjE,
      simp add:normal-valuation-belonging-to-prime-divisor-def,
      frule-tac  $v = v$  in valuation-some-in-prime-divisor,
      frule-tac  $v = \text{SOME } w. w \in P_K v$  in
        n-val-equiv-val,
      frule-tac  $v = v$  in some-in-P-equiv,
      frule-tac  $v = v$  and  $v' = \text{SOME } w. w \in P_K v$  and  $v'' =$ 
         $n\text{-val } K (\text{SOME } w. w \in P_K v)$  in val-equiv-axiom3,
      assumption+,
      frule-tac  $v = v$  in n-val-valuation,
      simp add:prime-divisor-def, simp add:n-val-valuation)
done

```

```

lemma (in Corps)  $\text{val-equiv-eq-pdiv}:\llbracket P \in Pds_K; P' \in Pds_K; \text{valuation } K v;$ 
   $\text{valuation } K v'; v\text{-equiv } K v v'; v \in P; v' \in P' \rrbracket \implies P = P'$ 
apply (simp add:prime-divisors-def,
      (erule exE)+, (erule conjE)+)
apply (rename-tac  $w w'$ ,
      frule-tac  $v = w$  in vals-in-P-equiv[of -  $v$ ], simp,
      frule-tac  $v = w'$  in vals-in-P-equiv[of -  $v'$ ], simp,
      frule-tac  $v = w$  and  $v' = v$  and  $v'' = v'$  in val-equiv-axiom3,
      assumption+,
      frule-tac  $v = w'$  in val-equiv-axiom2[of -  $v'$ ], assumption+,
      frule-tac  $v = w$  and  $v' = v'$  and  $v'' = w'$  in val-equiv-axiom3,
      assumption+) apply simp+
apply (simp add:P-eq-val-equiv)
done

```

```

lemma (in Corps)  $\text{distinct-p-divisors}:\llbracket P \in Pds_K; P' \in Pds_K \rrbracket \implies$ 
   $(\neg P = P') = (\neg v\text{-equiv } K (\nu_K P) (\nu_K P'))$ 
apply (rule iffI,
      rule contrapos-pp, simp+,
      frule val-equiv-eq-pdiv[of  $P P' \nu_K P \nu_K P'$ ], assumption+,

```



```

      simp add: representative-of-pd-valuation,
      simp add: representative-of-pd-valuation, assumption)
apply (rule n-val-representative[of P], assumption,
      rule n-val-representative[of P], assumption,
      simp,
      rule contrapos-pp, simp+, frule sym, thin-tac P = P',
      simp,
      frule representative-of-pd-valuation[of P],
      frule val-equiv-axiom1[of  $\nu_K P$ ], simp)
done

```

2.8 Approximation

definition

```

valuations :: [-, nat, nat  $\Rightarrow$  ('r  $\Rightarrow$  ant)]  $\Rightarrow$  bool where
valuations K n vv  $\longleftrightarrow$  ( $\forall j \leq n$ . valuation K (vv j))

```

definition

```

vals-nonequiv :: [-, nat, nat  $\Rightarrow$  ('r  $\Rightarrow$  ant)]  $\Rightarrow$  bool where
vals-nonequiv K n vv  $\longleftrightarrow$  valuations K n vv  $\wedge$ 
( $\forall j \leq n$ .  $\forall l \leq n$ .  $j \neq l \longrightarrow \neg$  (v-equiv K (vv j) (vv l)))

```

definition

```

Ostrowski-elim :: [-, nat, nat  $\Rightarrow$  ('b  $\Rightarrow$  ant), 'b]  $\Rightarrow$  bool where
Ostrowski-elim K n vv x  $\longleftrightarrow$ 
  ( $0 < (vv\ 0\ (1_r K \pm_K (-_a K\ x)))$ )  $\wedge$  ( $\forall j \in \text{nset}\ (Suc\ 0)\ n$ .  $0 < (vv\ j\ x)$ )

```

lemma (in Corps) Ostrowski-elim-0: $\llbracket \text{vals-nonequiv } K\ n\ vv; x \in \text{carrier } K; \text{Ostrowski-elim } K\ n\ vv\ x \rrbracket \Longrightarrow 0 < (vv\ 0\ (1_r \pm (-_a\ x)))$
apply (simp add: Ostrowski-elim-def)
done

lemma (in Corps) Ostrowski-elim-Suc: $\llbracket \text{vals-nonequiv } K\ n\ vv; x \in \text{carrier } K; \text{Ostrowski-elim } K\ n\ vv\ x; j \in \text{nset}\ (Suc\ 0)\ n \rrbracket \Longrightarrow 0 < (vv\ j\ x)$
apply (simp add: Ostrowski-elim-def)
done

lemma (in Corps) vals-nonequiv-valuation: $\llbracket \text{vals-nonequiv } K\ n\ vv; m \leq n \rrbracket \Longrightarrow \text{valuation } K\ (vv\ m)$
apply (simp add: vals-nonequiv-def, erule conjE)
apply (thin-tac $\forall j \leq n$. $\forall l \leq n$. $j \neq l \longrightarrow \neg$ v-equiv K (vv j) (vv l))
apply (simp add: valuations-def)
done

lemma (in Corps) vals-nonequiv: $\llbracket \text{vals-nonequiv } K\ (Suc\ (Suc\ n))\ vv; i \leq (Suc\ (Suc\ n)); j \leq (Suc\ (Suc\ n)); i \neq j \rrbracket \Longrightarrow \neg$ (v-equiv K (vv i) (vv j))

apply (*simp add:vals-nonequiv-def*)
done

lemma (*in Corps*) *skip-vals-nonequiv:vals-nonequiv* K (*Suc* (*Suc* n)) $vv \implies$
 $vals-nonequiv\ K\ (Suc\ n)\ (compose\ \{l.\ l \leq (Suc\ n)\}\ vv\ (skip\ j))$
apply (*subst vals-nonequiv-def*)
apply (*rule conjI*)
apply (*subst valuations-def, rule allI, rule impI,*
simp add:compose-def)
apply (*cut-tac l = ja and n = Suc n in skip-mem[of - - j], simp,*
frule-tac m = skip j ja in vals-nonequiv-valuation[of
Suc (Suc n) vv], simp, assumption)
apply (*(rule allI, rule impI)+, rule impI,*
cut-tac l = ja and n = Suc n in skip-mem[of - - j], simp,
cut-tac l = l and n = Suc n in skip-mem[of - - j], simp+)
apply (*cut-tac i = ja and j = l in skip-inj[of - Suc n - j], simp+,*
simp add:compose-def,
rule-tac i = skip j ja and j = skip j l in
vals-nonequiv[of n], assumption+)
done

lemma (*in Corps*) *not-v-equiv-reflex*: $\llbracket valuation\ K\ v; valuation\ K\ v';$
 $\neg v-equiv\ K\ v\ v' \rrbracket \implies \neg v-equiv\ K\ v'\ v$
apply (*simp add:v-equiv-def*)
done

lemma (*in Corps*) *nonequiv-ex-Ostrowski-elem*: $\llbracket valuation\ K\ v; valuation\ K\ v';$
 $\neg v-equiv\ K\ v\ v' \rrbracket \implies \exists x \in carrier\ K.\ 0 \leq (v\ x) \wedge (v'\ x) < 0$
apply (*subgoal-tac $\neg (\forall x \in carrier\ K.\ 0 \leq (v\ x) \longrightarrow 0 \leq (v'\ x))$*)
prefer 2
apply (*rule contrapos-pp, simp+,*
frule valuations-equiv[of v v'], assumption+,
simp add:val-equiv-axiom2[of v v'])
apply (*simp, erule bexE, erule conjE, simp add:aneg-le*)
apply *blast*
done

lemma (*in Corps*) *field-op-minus*: $\llbracket a \in carrier\ K; b \in carrier\ K; b \neq 0 \rrbracket \implies$
 $-_a\ (a \cdot_r\ (b^{-K})) = (-_a\ a) \cdot_r\ (b^{-K})$
apply (*cut-tac invf-closed1[of b], simp,*
erule conjE, cut-tac field-is-ring,
simp add:Ring.ring-inv1[of K a b^{-K}], simp)
done

lemma (*in Corps*) *field-one-plus-frac1*: $\llbracket a \in carrier\ K; b \in carrier\ K; b \neq 0 \rrbracket$
 $\implies 1_r \pm (a \cdot_r\ (b^{-K})) = (b \pm a) \cdot_r\ (b^{-K})$
apply (*cut-tac field-is-ring,*
cut-tac invf-closed1[of b], simp+, erule conjE,
cut-tac field-is-idom)

apply (rule *Idomain.idom-mult-cancel-r*[of K $1_r \pm (a \cdot_r (b^{-K}))$]
 $(b \pm a) \cdot_r (b^{-K})$ b], assumption+,
 frule *Idomain.idom-is-ring*[of K], frule *Ring.ring-is-ag*[of K],
 rule *aGroup.ag-pOp-closed* [of K], assumption+,
 simp add: *Ring.ring-one*, rule *Ring.ring-tOp-closed*, assumption+,
 rule *Ring.ring-tOp-closed*, assumption+,
 frule *Ring.ring-is-ag*[of K],
 rule *aGroup.ag-pOp-closed*, assumption+,
 subst *Ring.ring-distrib2*[of K b], assumption+,
 simp add: *Ring.ring-one*, simp add: *Ring.ring-tOp-closed*,
 simp add: *Ring.ring-l-one*) **thm** *Ring.ring-distrib2*[of K b^{-K}]
apply (subst *Ring.ring-distrib2*[of K b^{-K}], assumption+,
 simp add: *Ring.ring-tOp-commute*[of K b b^{-K}],
 subst *linvf*[of b], simp,
 subst *Ring.ring-distrib2*[of K b], assumption+,
 simp add: *Ring.ring-one*, simp add: *Ring.ring-tOp-closed*,
 simp add: *Ring.ring-l-one*, simp)
done

lemma (in *Corps*) *field-one-plus-frac2*: $\llbracket a \in \text{carrier } K; b \in \text{carrier } K;$
 $a \pm b \neq \mathbf{0} \rrbracket \implies 1_r \pm (-_a (a \cdot_r (a \pm b)^{-K})) = b \cdot_r ((a \pm b)^{-K})$
apply (frule *field-op-minus*[of a $a \pm b$],
 cut-tac *field-is-ring*, frule *Ring.ring-is-ag*[of K],
 simp add: *aGroup.ag-pOp-closed*, assumption, simp,
 thin-tac $-_a (a \cdot_r (a \pm b)^{-K}) = (-_a a) \cdot_r (a \pm b)^{-K}$)
apply (cut-tac *field-is-ring*, frule *Ring.ring-is-ag*[of K],
 frule *aGroup.ag-mOp-closed*[of K a], assumption,
 frule *field-one-plus-frac1*[of $-_a a$ $a \pm b$],
 simp add: *aGroup.ag-pOp-closed*, simp, simp,
 thin-tac $1_r \pm (-_a a) \cdot_r (a \pm b)^{-K} = (a \pm b \pm -_a a) \cdot_r (a \pm b)^{-K}$,
 simp add: *aGroup.ag-pOp-assoc*[of K a b $-_a a$],
 simp add: *aGroup.ag-pOp-commute*[of K b $-_a a$],
 simp add: *aGroup.ag-pOp-assoc*[*THEN sym*],
 simp add: *aGroup.ag-r-inv1*,
 simp add: *aGroup.ag-l-zero*)
done

lemma (in *Corps*) *field-one-plus-frac3*: $\llbracket x \in \text{carrier } K; x \neq 1_r;$
 $1_r \pm x \cdot_r (1_r \pm -_a x) \neq \mathbf{0} \rrbracket \implies$
 $1_r \pm -_a x \cdot_r (1_r \pm x \cdot_r (1_r \pm -_a x))^{-K} =$
 $(1_r \pm -_a x \cdot_r (1_r \pm x \cdot_r (1_r \pm -_a x))^{-K} (1_r \pm -_a x \cdot_r (1_r \pm x \cdot_r (1_r \pm -_a x))^{-K})) \cdot_r (1_r \pm x \cdot_r (1_r \pm -_a x))^{-K}$
apply (cut-tac *field-is-ring*, frule *Ring.ring-is-ag*, frule *Ring.ring-one*,
 cut-tac *invf-closed1*[of $1_r \pm x \cdot_r (1_r \pm -_a x)$], simp, erule *conjE*)
apply (subst *Ring.ring-inv1-1*, assumption+,
 subst *field-one-plus-frac1*[of $-_a x$ $1_r \pm x \cdot_r (1_r \pm -_a x)$])
apply (rule *aGroup.ag-mOp-closed*, assumption+,
 rule *aGroup.ag-pOp-closed*, assumption+,
 rule *Ring.ring-tOp-closed*, assumption+)
apply (rule *aGroup.ag-pOp-closed*, assumption+, rule *aGroup.ag-mOp-closed*,

```

    assumption+,
    subst Ring.ring-distrib1, assumption+,
    rule aGroup.ag-mOp-closed, assumption+)
apply (simp add:Ring.ring-r-one)
apply (simp add:Ring.ring-inv1-2[THEN sym, of K x x])
apply (subgoal-tac  $1_r \pm (x \pm -_a x \cdot_r x) \pm -_a x = 1_r \pm -_a x \sim^K (Suc (Suc 0))$ ,
    simp,
    frule Ring.ring-tOp-closed[of K x x], assumption+)

apply (frule Ring.ring-tOp-closed[of K x x], assumption+,
    frule aGroup.ag-mOp-closed[of K x  $\cdot_r$  x], assumption+,
    frule aGroup.ag-mOp-closed[of K x], assumption+)
apply (subst aGroup.ag-pOp-assoc, assumption+,
    rule aGroup.ag-pOp-closed, assumption+)
apply (rule aGroup.ag-pOp-add-l[of K  $x \pm -_a x \cdot_r x \pm -_a x$ 
     $-_a x \sim^K (Suc (Suc 0))$   $1_r$ ], assumption+,
    (rule aGroup.ag-pOp-closed, assumption+)+,
    rule aGroup.ag-mOp-closed, assumption+, rule Ring.npClose,
    assumption+,
    subst aGroup.ag-pOp-commute, assumption+,
    simp add:aGroup.ag-pOp-assoc aGroup.ag-r-inv1 aGroup.ag-r-zero)
apply (simp add:Ring.ring-l-one)
apply simp
apply (rule aGroup.ag-pOp-closed, assumption+,
    rule Ring.ring-tOp-closed, assumption+,
    rule aGroup.ag-pOp-closed, assumption+,
    rule aGroup.ag-mOp-closed[of K x], assumption+)
done

lemma (in Corps) OstrowskiTr1:  $\llbracket \text{valuation } K \ v; s \in \text{carrier } K; t \in \text{carrier } K; \$ 
 $0 \leq (v \ s); v \ t < 0 \rrbracket \implies s \pm t \neq \mathbf{0}$ 
apply (rule contrapos-pp, simp+,
    cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    simp only:aGroup.ag-plus-zero[THEN sym, of K s t])
apply (simp add:val-minus-eq[of v t])
done

lemma (in Corps) OstrowskiTr2:  $\llbracket \text{valuation } K \ v; s \in \text{carrier } K; t \in \text{carrier } K; \$ 
 $0 \leq (v \ s); v \ t < 0 \rrbracket \implies 0 < (v \ (1_r \pm (-_a ((t \cdot_r ((s \pm t)^{-K}))))))$ 
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    frule OstrowskiTr1[of v s t], assumption+,
    frule field-one-plus-frac2[of t s], assumption+,
    simp add:aGroup.ag-pOp-commute)
apply (subst aGroup.ag-pOp-commute[of K s t], assumption+, simp,
    simp add:aGroup.ag-pOp-commute[of K t s],
    thin-tac  $1_r \pm -_a (t \cdot_r (s \pm t)^{-K}) = s \cdot_r (s \pm t)^{-K}$ ,
    frule aGroup.ag-pOp-closed[of K s t], assumption+,
    cut-tac invf-closed1[of s  $\pm$  t], simp, erule conjE)
apply (simp add:val-t2p,

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```

    simp add:value-of-inv,
    frule aless-le-trans[of v t 0 v s], assumption+,
    frule value-less-eq[THEN sym, of v t s], assumption+,
    simp add:aGroup.ag-pOp-commute,
    frule aless-diff-poss[of v t v s], simp add:diff-ant-def, simp)
done

lemma (in Corps) OstrowskiTr3:[[valuation K v; s ∈ carrier K; t ∈ carrier K;
  0 ≤ (v t); v s < 0]] ⇒ 0 < (v (t ·r ((s ± t)-K)))
apply (frule aless-le-trans[of v s 0 v t], assumption+,
  cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
  frule aGroup.ag-pOp-closed[of K s t], assumption+,
  frule OstrowskiTr1[of v t s], assumption+,
  frule value-less-eq[THEN sym, of v s t], assumption+)
apply (simp add:aGroup.ag-pOp-commute[of K t s],
  cut-tac invf-closed1[of s ± t], simp) apply (
  erule conjE, simp add:val-t2p[of v], simp add:value-of-inv)
apply (cut-tac aless-diff-poss[of v s v t],
  simp add:diff-ant-def, simp+)
done

lemma (in Corps) restrict-Ostrowski-elem:[[x ∈ carrier K;
  Ostrowski-elem K (Suc (Suc n)) vv x]] ⇒ Ostrowski-elem K (Suc n) vv x
apply (simp add:Ostrowski-elem-def,
  erule conjE, rule ballI, simp add:nset-def,
  insert lessI [of Suc n])
done

lemma (in Corps) restrict-vals-nonequiv:vals-nonequiv K (Suc (Suc n)) vv ⇒
  vals-nonequiv K (Suc n) vv
apply (simp add:vals-nonequiv-def,
  erule conjE, simp add:valuations-def)
done

lemma (in Corps) restrict-vals-nonequiv1:vals-nonequiv K (Suc (Suc n)) vv ⇒
  vals-nonequiv K (Suc n) (compose {h. h ≤ (Suc n)} vv (skip 1))
apply (simp add:vals-nonequiv-def, (erule conjE)+,
  rule conjI,
  thin-tac ∀j≤Suc (Suc n). ∀l≤Suc (Suc n). j ≠ l ⇒
    ¬ v-equiv K (vv j) (vv l),
  simp add:valuations-def, rule allI, rule impI,
  simp add:compose-def skip-def nset-def)
apply ((rule allI, rule impI)+, rule impI)
apply (simp add:compose-def skip-def nset-def)
done

lemma (in Corps) restrict-vals-nonequiv2:[[vals-nonequiv K (Suc (Suc n)) vv]
  ⇒ vals-nonequiv K (Suc n) (compose {j. j ≤ (Suc n)} vv (skip 2))
apply (simp add:vals-nonequiv-def, (erule conjE)+,

```

rule conjI,
 $\text{thin-tac } \forall j \leq \text{Suc } (\text{Suc } n). \forall l \leq \text{Suc } (\text{Suc } n). j \neq l \longrightarrow$
 $\qquad \neg v\text{-equiv } K (vv\ j) (vv\ l),$
 $\text{simp add:valuations-def,}$
 $\text{rule allI, rule impI)}$
apply (simp add:compose-def skip-def nset-def,
(rule allI, rule impI)+, rule impI,
simp add:compose-def skip-def nset-def)
done

lemma (in Corps) OstrowskiTr31: $\llbracket \text{valuation } K\ v; s \in \text{carrier } K;$
 $0 < (v\ (1_r \pm (-_a\ s))) \rrbracket \implies s \neq \mathbf{0}$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (rule contrapos-pp, simp+)
apply (simp add:aGroup.ag-inv-zero,
frule Ring.ring-one[of K], simp add:aGroup.ag-r-zero)
apply (simp add:value-of-one)
done

lemma (in Corps) OstrowskiTr32: $\llbracket \text{valuation } K\ v; s \in \text{carrier } K;$
 $0 < (v\ (1_r \pm (-_a\ s))) \rrbracket \implies 0 \leq (v\ s)$
apply (rule contrapos-pp, simp+,
cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
simp add:aneg-le,
frule has-val-one-neg-zero[of v])
apply (frule OstrowskiTr31[of v s], assumption+,
frule not-sym,
frule Ring.ring-one[of K])
apply (frule value-less-eq[THEN sym, of v -_a s 1_r],
simp add:aGroup.ag-mOp-closed, assumption+,
simp add:val-minus-eq)
apply (simp add:value-of-one,
frule aGroup.ag-mOp-closed[of K s], assumption+,
simp add:aGroup.ag-pOp-commute[of K -_a s 1_r],
simp add:val-minus-eq)
done

lemma (in Corps) OstrowskiTr4: $\llbracket \text{valuation } K\ v; s \in \text{carrier } K; t \in \text{carrier } K;$
 $0 < (v\ (1_r \pm (-_a\ s))) ; 0 < (v\ (1_r \pm (-_a\ t))) \rrbracket \implies$
 $0 < (v\ (1_r \pm (-_a\ (s \cdot_r t))))$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
frule Ring.ring-one[of K])
apply (subgoal-tac $1_r \pm (-_a\ (s \cdot_r t)) =$
 $1_r \pm (-_a\ s) \pm (s \cdot_r (1_r \pm (-_a\ t))), \text{simp,}$
thin-tac $1_r \pm (-_a\ (s \cdot_r t)) = 1_r \pm -_a\ s \pm s \cdot_r (1_r \pm -_a\ t))$
apply (frule aGroup.ag-mOp-closed[of K s], assumption+,
frule aGroup.ag-pOp-closed[of K 1_r -_a s], assumption+,
frule aGroup.ag-mOp-closed[of K t], assumption+,
frule aGroup.ag-pOp-closed[of K 1_r -_a t], assumption+),

```

    frule Ring.ring-tOp-closed[of K s 1r ± (-a t)], assumption+,
    frule amin-le-plus[of v 1r ± (-a s) s ·r (1r ± (-a t))], assumption+
  apply (frule amin-gt[of 0 v (1r ± -a s) v (s ·r (1r ± -a t))])
  apply (simp add:val-t2p,
    frule OstrowskiTr32[of v s], assumption+,
    rule aadd-pos-poss[of v s v (1r ± -a t)], assumption+,
    simp add:Ring.ring-distrib1)
  apply (frule aGroup.ag-mOp-closed[of K t], assumption,
    simp add:Ring.ring-distrib1 Ring.ring-r-one,
    frule aGroup.ag-mOp-closed[of K s], assumption+,
    subst aGroup.pOp-assocTr43, assumption+,
    simp add:Ring.ring-tOp-closed,
    simp add:aGroup.ag-l-inv1 aGroup.ag-r-zero,
    simp add:Ring.ring-inv1-2)
done

```

```

lemma (in Corps) OstrowskiTr5: [| vals-nonequiv K (Suc (Suc n)) vv;
  s ∈ carrier K; t ∈ carrier K;
  0 ≤ (vv (Suc 0)) s ∧ 0 ≤ (vv (Suc (Suc 0))) t;
  Ostrowski-elem K (Suc n) (compose {j. j ≤ (Suc n)} vv (skip 1)) s;
  Ostrowski-elem K (Suc n) (compose {j. j ≤ (Suc n)} vv (skip 2)) t |] ⇒
  Ostrowski-elem K (Suc (Suc n)) vv (s ·r t)
  apply (erule conjE,
    cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    frule-tac x = s and y = t in Ring.ring-tOp-closed[of K], assumption+,
    frule skip-vals-nonequiv[of n vv 1],
    frule skip-vals-nonequiv[of n vv 2],
    subst Ostrowski-elem-def, rule conjI)

```

```

  apply (rule OstrowskiTr4,
    simp add:vals-nonequiv-valuation[of Suc (Suc n) vv 0],
    assumption+,
    frule Ostrowski-elem-0[of Suc n
      compose {j. j ≤ (Suc n)} vv (skip 1) s], assumption+,
    simp add:skip-def compose-def,
    frule Ostrowski-elem-0[of Suc n
      compose {j. j ≤ (Suc n)} vv (skip 2) t], assumption+,
    simp add:skip-def compose-def)

```

```

  apply (rule ballI,
    case-tac j = Suc 0,
    frule-tac j = Suc 0 in Ostrowski-elem-Suc[of Suc n
      compose {j. j ≤ (Suc n)} vv (skip 2) t], assumption+,
    simp add:nset-def) apply (
    thin-tac Ostrowski-elem K (Suc n) (compose {j. j ≤ Suc n} vv (skip 1)) s,
    thin-tac Ostrowski-elem K (Suc n) (compose {j. j ≤ Suc n} vv (skip 2)) t,
    thin-tac vals-nonequiv K (Suc n) (compose {l. l ≤ Suc n} vv (skip 1)),
    frule vals-nonequiv-valuation[of Suc n
      compose {j. j ≤ (Suc n)} vv (skip 2) Suc 0])

```

apply *simp+*
apply (*simp add:skip-def compose-def*,
simp add:val-t2p, simp add:aadd-pos-poss)

apply (*case-tac j = Suc (Suc 0)*,
frule vals-nonequiv-valuation[of Suc n
compose {j. j ≤ Suc n} vv (skip 1) Suc 0],
simp,
frule-tac j = Suc 0 in Ostrowski-elem-Suc[of Suc n
compose {j. j ≤ (Suc n)} vv (skip 1) s],
assumption+, *simp add:nset-def*,
simp add:skip-def compose-def,
simp add:val-t2p, rule aadd-poss-pos, assumption+)
apply (*frule-tac j = j in nsetTr1[of - Suc 0 Suc (Suc n)]*, *assumption*,
frule-tac j = j in nsetTr2[of - Suc 0 Suc n],
thin-tac j ∈ nset (Suc (Suc 0)) (Suc (Suc n)),
frule-tac j = j - Suc 0 in Ostrowski-elem-Suc[of Suc n
compose {j. j ≤ (Suc n)} vv (skip 1) s], *assumption+*)

apply (*frule-tac j = j - Suc 0 in Ostrowski-elem-Suc[of Suc n*
compose {j. j ≤ (Suc n)} vv (skip 2) t], *assumption+*,
thin-tac Ostrowski-elem K (Suc n) (compose {j. j ≤ (Suc n)} vv (skip 1)) s,
thin-tac Ostrowski-elem K (Suc n) (compose {j. j ≤ (Suc n)} vv (skip 2)) t,
thin-tac vals-nonequiv K (Suc n) (compose {j. j ≤ (Suc n)} vv (skip 1)),
thin-tac vals-nonequiv K (Suc n) (compose {j. j ≤ (Suc n)} vv (skip 2)))

apply (*simp add:compose-def skip-def nset-def*,
(erule conjE)+, simp, subgoal-tac ¬ (j - Suc 0 ≤ Suc 0), simp)
apply (*frule-tac m = j in vals-nonequiv-valuation[of Suc (Suc n)]*,
assumption+,
simp add:val-t2p,
rule-tac x = vv j s and y = vv j t in aadd-pos-poss,
simp add:aless-imp-le, assumption)

apply *simp*
done

lemma (*in Corps*) *one-plus-x-nonzero*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; v \ x < 0 \rrbracket$
 $\implies 1_r \pm x \in \text{carrier } K \wedge v \ (1_r \pm x) < 0$

apply (*cut-tac field-is-ring, frule Ring.ring-is-ag[of K]*,
frule Ring.ring-one[of K],
frule aGroup.ag-pOp-closed[of K 1_r x], *assumption+*,
simp)

apply (*frule value-less-eq[of v x 1_r]*, *assumption+*,
simp add:value-of-one, simp add:aGroup.ag-pOp-commute)

done

lemma (*in Corps*) *val-neg-nonzero*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; v \ x < 0 \rrbracket \implies$
 $x \neq 0$

apply (rule *contrapos-pp*, *simp+*, *simp add:value-of-zero*)
apply (frule *aless-imp-le*[of ∞ 0],
 cut-tac *inf-ge-any*[of 0],
 frule *ale-antisym*[of 0 ∞], *assumption+*, *simp*)
done

lemma (in *Corps*) *OstrowskiTr6*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; \neg 0 \leq (v \ x) \rrbracket \implies$
 $(1_r \pm x \cdot_r (1_r \pm -_a x)) \in \text{carrier } K - \{0\}$
apply (*simp add:aneg-le*,
 cut-tac *field-is-ring*, frule *Ring.ring-is-ag*[of K],
 frule *aGroup.ag-mOp-closed*[of $K \ x$], *assumption+*,
 frule *one-plus-x-nonzero*[of $v \ -_a \ x$], *assumption+*,
simp add:val-minus-eq, *erule conjE*)

apply (rule *conjI*,
 rule *aGroup.ag-pOp-closed*[of K], *assumption+*,
simp add:Ring.ring-one, rule *Ring.ring-tOp-closed*[of K], *assumption+*)

apply (frule *val-t2p*[of $v \ x \ 1_r \pm (-_a \ x)$], *assumption+*,
 frule *val-neg-nonzero*[of $v \ x$], *assumption+*,
 frule *val-nonzero-z*[of $v \ x$], *assumption+*, *erule exE*,
 frule-tac $z = z$ **in** *aadd-less-mono-z*[of $v \ (1_r \pm (-_a \ x)) \ 0$],
simp add:aadd-0-l,
simp only:aadd-commute[of $v \ (1_r \pm -_a \ x)$],
 frule-tac $x = \text{ant } z + v \ (1_r \pm -_a \ x)$ **and** $y = \text{ant } z$ **in**
aless-trans[of $- \ - \ 0$], *assumption*,
drule sym, *simp*)

apply (frule-tac $x = x$ **and** $y = 1_r \pm -_a \ x$ **in** *Ring.ring-tOp-closed*[of K],
assumption+,
 frule *one-plus-x-nonzero*[of $v \ x \cdot_r (1_r \pm (-_a \ x))$],
assumption+, *erule conjE*,
 rule *val-neg-nonzero*[of v], *assumption+*)
done

lemma (in *Corps*) *OstrowskiTr7*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; \neg 0 \leq (v \ x) \rrbracket \implies$
 $1_r \pm -_a (x \cdot_r ((1_r \pm x \cdot_r (1_r \pm -_a x))^{-K})) =$
 $(1_r \pm -_a x \pm x \cdot_r (1_r \pm -_a x)) \cdot_r ((1_r \pm x \cdot_r (1_r \pm -_a x))^{-K})$
apply (*cut-tac field-is-ring*,
 frule *OstrowskiTr6*[of $v \ x$], *assumption+*, *simp*, *erule conjE*,
 cut-tac *field-is-idom*,
 cut-tac *invf-closed1*[of $1_r \pm x \cdot_r (1_r \pm -_a x)$], *simp*,
 frule *Ring.ring-is-ag*[of K],
 frule *aGroup.ag-mOp-closed*[of $K \ x$], *assumption+*,
 frule *Ring.ring-one*[of K],
 frule *aGroup.ag-pOp-closed*[of $K \ 1_r -_a \ x$], *assumption+*,
 rule *Idomain.idom-mult-cancel-r*[of $K \ 1_r \pm -_a (x \cdot_r ((1_r \pm x \cdot_r$
 $(1_r \pm -_a x))^{-K})) \ (1_r \pm -_a x \pm x \cdot_r (1_r \pm -_a x)) \cdot_r$
 $((1_r \pm x \cdot_r (1_r \pm -_a x))^{-K}) \ (1_r \pm x \cdot_r (1_r \pm -_a x))$],

```

    assumption+)
apply (rule aGroup.ag-pOp-closed, assumption+, rule aGroup.ag-mOp-closed,
    assumption+,
    rule Ring.ring-tOp-closed, assumption+, simp, rule Ring.ring-tOp-closed,
    assumption+,
    (rule aGroup.ag-pOp-closed, assumption+)+,
    rule Ring.ring-tOp-closed, assumption+, simp, assumption+,
    subst Ring.ring-tOp-assoc, assumption+,
    rule aGroup.ag-pOp-closed, assumption+,
    simp add:Ring.ring-tOp-closed, simp, simp)
apply (subst linvf[of  $1_r \pm x \cdot_r (1_r \pm -_a x)$ ], simp,
    (subst Ring.ring-distrib2, assumption+)+, erule conjE)
apply (rule aGroup.ag-mOp-closed, assumption,
    rule Ring.ring-tOp-closed, assumption+,
    subst Ring.ring-r-one, assumption+)
apply (rule aGroup.ag-pOp-closed, assumption+,
    rule Ring.ring-tOp-closed, assumption+,
    erule conjE,
    simp add:Ring.ring-inv1-1,
    simp add:Ring.ring-tOp-assoc[of  $K -_a x (1_r \pm x \cdot_r (1_r \pm -_a x))^{-K}$ ],
    simp add:linvf, simp add:Ring.ring-r-one Ring.ring-l-one,
    frule Ring.ring-tOp-closed[of  $K x 1_r \pm -_a x$ ], assumption+,
    simp add:aGroup.ag-pOp-assoc, simp add:aGroup.ag-pOp-commute)
apply simp
done

lemma (in Corps) Ostrowski-elem-nonzero:  $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv; \\ x \in \text{carrier } K; \text{Ostrowski-elem } K \text{ (Suc } n) \text{ } vv \ x \rrbracket \implies x \neq \mathbf{0}$ 
apply (simp add:Ostrowski-elem-def,
    frule conjunct1, fold Ostrowski-elem-def,
    frule vals-nonequiv-valuation[of Suc n vv 0], simp)
apply (rule contrapos-pp, simp+,
    cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    simp add:aGroup.ag-inv-zero, frule Ring.ring-one[of K],
    simp add:aGroup.ag-r-zero, simp add:value-of-one)
done

lemma (in Corps) Ostrowski-elem-not-one:  $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv; \\ x \in \text{carrier } K; \text{Ostrowski-elem } K \text{ (Suc } n) \text{ } vv \ x \rrbracket \implies 1_r \pm -_a x \neq \mathbf{0}$ 
apply (frule vals-nonequiv-valuation [of Suc n vv Suc 0],
    simp,
    simp add:Ostrowski-elem-def, frule conjunct2,
    fold Ostrowski-elem-def)
apply (subgoal-tac  $0 < (vv \text{ (Suc } 0) \ x)$ ,
    rule contrapos-pp, simp+,
    cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    frule Ring.ring-one[of K],
    simp only:aGroup.ag-eq-diffzero[THEN sym, of  $K 1_r x$ ],
    drule sym, simp, simp add:value-of-one,

```

```

      subgoal-tac Suc 0 ∈ nset (Suc 0) (Suc n), simp,
      simp add:nset-def)
done

lemma (in Corps) val-unit-cond: [ valuation K v; x ∈ carrier K;
  0 < (v (1r ± -a x)) ] ⇒ v x = 0
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
  frule Ring.ring-one[of K])

apply (frule aGroup.ag-mOp-closed[of K 1r], assumption+,
  frule has-val-one-neq-zero[of v])

apply (frule aGroup.ag-pOp-assoc[of K -a 1r 1r -a x], assumption+,
  simp add:aGroup.ag-mOp-closed, simp add:aGroup.ag-l-inv1,
  frule aGroup.ag-mOp-closed[of K x], assumption+,
  simp add:aGroup.ag-l-zero)
apply (subgoal-tac v (-a x) = v (-a 1r ± (1r ± -a x))) prefer 2
  apply simp
apply (thin-tac -a x = -a 1r ± (1r ± -a x),
  frule value-less-eq[of v -a 1r 1r ± -a x],
  assumption+,
  rule aGroup.ag-pOp-closed, assumption+,
  simp add:val-minus-eq value-of-one, simp add:val-minus-eq)
apply (simp add: value-of-one)
done

end

theory Valuation2
imports Valuation1
begin

lemma (in Corps) OstrowskiTr8: [ valuation K v; x ∈ carrier K;
  0 < v (1r ± -a x) ] ⇒
  0 < (v (1r ± -a (x ·r ((1r ± x ·r (1r ± -a x))-K))))
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (frule aGroup.ag-mOp-closed[of K x], assumption+,
  frule Ring.ring-one[of K],
  frule aGroup.ag-pOp-closed[of K 1r -a x], assumption+,
  frule OstrowskiTr32[of v x], assumption+)
apply (case-tac x = 1r K, simp,
  simp add:aGroup.ag-r-inv1, simp add:Ring.ring-times-x-0,
  simp add:aGroup.ag-r-zero, cut-tac val-field-1-neq-0,
  cut-tac invf-one, simp, simp add:Ring.ring-r-one,
  simp add:aGroup.ag-r-inv1, assumption+)
apply (frule aGroup.ag-pOp-closed[of K 1r x ·r (1r ± -a x)], assumption+,
  rule Ring.ring-tOp-closed, assumption+)
apply (cut-tac invf-closed[of 1r ± x ·r (1r ± -a x)])

```

apply (*cut-tac field-one-plus-frac3*[of x], *simp del:npow-suc*,
subst val-t2p[of v], *assumption*+)

apply (*rule aGroup.ag-pOp-closed*, *assumption*+,
rule aGroup.ag-mOp-closed, *assumption*+, *rule Ring.npClose*,
assumption+,
thin-tac $1_r \pm -_a x \cdot_r (1_r \pm x \cdot_r (1_r \pm -_a x))^{-K} =$
 $(1_r \pm -_a x^{-K} (\text{Suc} (\text{Suc } 0))) \cdot_r (1_r \pm x \cdot_r (1_r \pm -_a x))^{-K}$,
subgoal-tac $1_r \pm -_a x^{-K} (\text{Suc} (\text{Suc } 0)) = (1_r \pm x) \cdot_r (1_r \pm -_a x)$,
simp del:npow-suc,
thin-tac $1_r \pm -_a x^{-K} (\text{Suc} (\text{Suc } 0)) = (1_r \pm x) \cdot_r (1_r \pm -_a x)$)

apply (*subst val-t2p*[of v], *assumption*+,
rule aGroup.ag-pOp-closed, *assumption*+,
subst value-of-inv[of $v \ 1_r \pm x \cdot_r (1_r \pm -_a x)$], *tactic* $\langle \text{CHANGED distinct-subgoals-tac} \rangle$, *assumption*+)

apply (*rule contrapos-pp*, *simp*+,
frule Ring.ring-tOp-closed[of $K \ x \ (1_r \pm -_a x)$], *assumption*+,
simp add:aGroup.ag-pOp-commute[of $K \ 1_r$],
frule aGroup.ag-eq-diffzero[*THEN sym*, of $K \ x \cdot_r (-_a x \pm 1_r) -_a 1_r$],
assumption+, *rule aGroup.ag-mOp-closed*, *assumption*+)

apply (*simp add:aGroup.ag-inv-inv*[of $K \ 1_r$],
frule eq-elems-eq-val[of $x \cdot_r (-_a x \pm 1_r) -_a 1_r \ v$],
thin-tac $x \cdot_r (-_a x \pm 1_r) = -_a 1_r$,
simp add:val-minus-eq value-of-one)

apply (*simp add:val-t2p*,
frule aadd-pos-poss[of $v \ x \ v \ (-_a x \pm 1_r)$], *assumption*+,
simp,
subst value-less-eq[*THEN sym*, of $v \ 1_r \ x \cdot_r (1_r \pm -_a x)$],
assumption+,
rule Ring.ring-tOp-closed, *assumption*+,
simp add:value-of-one, *subst val-t2p*[of v], *assumption*+,
rule aadd-pos-poss[of $v \ x \ v \ (1_r \pm -_a x)$], *assumption*+,
simp add:value-of-one,
cut-tac aadd-pos-poss[of $v \ (1_r \pm x) \ v \ (1_r \pm -_a x)$],
simp add:aadd-0-r, *rule val-axiom4*, *assumption*+)

apply (*subst Ring.ring-distrib2*, *assumption*+, *simp add:Ring.ring-l-one*,
subst Ring.ring-distrib1, *assumption*+, *simp add:Ring.ring-r-one*,
subst aGroup.pOp-assocTr43, *assumption*+,
rule Ring.ring-tOp-closed, *assumption*+,
simp add:aGroup.ag-l-inv1 aGroup.ag-r-zero,
subst Ring.ring-inv1-2, *assumption*+, *simp*, *assumption*+)

apply *simp*

apply (*rule contrapos-pp*, *simp*+,
frule Ring.ring-tOp-closed[of $K \ x \ (1_r \pm -_a x)$], *assumption*+,
simp add:aGroup.ag-pOp-commute[of $K \ 1_r$],
frule aGroup.ag-eq-diffzero[*THEN sym*, of $K \ x \cdot_r (-_a x \pm 1_r) -_a 1_r$],

```

    assumption+, rule aGroup.ag-mOp-closed, assumption+)
apply (simp add:aGroup.ag-inv-inv[of K 1r],
    frule eq-elems-eq-val[of x ·r (−a x ± 1r) −a 1r v],
    thin-tac x ·r (−a x ± 1r) = −a 1r,
    simp add:val-minus-eq value-of-one,
    simp add:val-t2p,
    frule aadd-pos-poss[of v x v (−a x ± 1r)], assumption+,
    simp)
done

lemma (in Corps) OstrowskiTr9:⟦valuation K v; x ∈ carrier K; 0 < (v x)⟧ ⇒
    0 < (v (x ·r ((1r ± x ·r (1r ± −a x))−K)))
apply (subgoal-tac 1r ± x ·r (1r ± −a x) ≠ 0)
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    frule aGroup.ag-mOp-closed[of K x], assumption+,
    frule Ring.ring-one[of K],
    frule aGroup.ag-pOp-closed[of K 1r −a x], assumption+,
    subst val-t2p, assumption+,
    cut-tac invf-closed1[of 1r ± x ·r (1r ± −a x)], simp)
apply simp
apply (rule aGroup.ag-pOp-closed, assumption+,
    rule Ring.ring-tOp-closed, assumption+)

apply (subst value-of-inv[of v 1r ± x ·r (1r ± −a x)], assumption+,
    rule aGroup.ag-pOp-closed, assumption+,
    rule Ring.ring-tOp-closed, assumption+,
    frule value-less-eq[THEN sym, of v 1r −a x], assumption+,
    simp add:value-of-one, simp add:val-minus-eq,
    subst value-less-eq[THEN sym, of v 1r x ·r (1r ± −a x)],
    assumption+, rule Ring.ring-tOp-closed, assumption+,
    simp add:value-of-one, subst val-t2p, assumption+,
    subst aadd-commute,
    rule aadd-pos-poss[of v (1r ± −a x) v x],
    simp, assumption, simp add:value-of-one,
    simp add:aadd-0-r)
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    frule Ring.ring-one,
    rule contrapos-pp, simp+,
    frule Ring.ring-tOp-closed[of K x (1r ± −a x)], assumption+,
    rule aGroup.ag-pOp-closed, assumption+,
    rule aGroup.ag-mOp-closed, assumption+,
    frule aGroup.ag-mOp-closed[of K x], assumption+)

apply (simp add:aGroup.ag-pOp-commute[of K 1r],
    frule aGroup.ag-eq-diffzero[THEN sym, of K x ·r (−a x ± 1r) −a 1r],
    simp add:aGroup.ag-pOp-commute,
    rule aGroup.ag-mOp-closed, assumption+,
    simp add:aGroup.ag-inv-inv,
    frule eq-elems-eq-val[of x ·r (−a x ± 1r) −a 1r v],

```

```

    thin-tac x ·r (-a x ± 1r) = -a 1r,
    simp add:val-minus-eq value-of-one,
    frule-tac aGroup.ag-pOp-closed[of K -a x 1r], assumption+,
    simp add:val-t2p)
  apply (simp add:aadd-commute[of v x],
    cut-tac aadd-pos-poss[of v (-a x ± 1r) v x], simp,
    subst aGroup.ag-pOp-commute, assumption+,
    subst value-less-eq[THEN sym, of v 1r -a x], assumption+,
    simp add:value-of-one val-minus-eq, simp add:value-of-one)

  apply assumption
done

lemma (in Corps) OstrowskiTr10:[[valuation K v; x ∈ carrier K;
  ¬ 0 ≤ v x]] ⇒ 0 < (v (x ·r ((1r ± x ·r (1r ± -a x))-K)))
  apply (frule OstrowskiTr6[of v x], assumption+,
    cut-tac invf-closed1[of 1r ± x ·r (1r ± -a x)], simp,
    erule conjE, simp add:aneg-le, frule val-neg-nonzero[of v x],
    (erule conjE)+, assumption+, erule conjE)
  apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    frule aGroup.ag-mOp-closed[of K x], assumption+,
    frule Ring.ring-one[of K],
    frule aGroup.ag-pOp-closed[of K 1r -a x], assumption+,
    subst val-t2p, assumption+)
  apply (subst value-of-inv[of v 1r ± x ·r (1r ± -a x)],
    assumption+, subst aGroup.ag-pOp-commute[of K 1r], assumption+,
    rule Ring.ring-tOp-closed, assumption+,
    subst value-less-eq[THEN sym, of v
      x ·r (1r ± -a x) 1r], assumption+)
  apply (rule Ring.ring-tOp-closed, assumption+, simp add:value-of-one,
    frule one-plus-x-nonzero[of v -a x], assumption,
    simp add:val-minus-eq, erule conjE, simp,
    subst val-t2p[of v], assumption+, simp add:aadd-two-negg)

  apply (simp add:val-t2p,
    frule value-less-eq[THEN sym, of v -a x 1r], assumption+,
    simp add:value-of-one, simp add:val-minus-eq,
    simp add:val-minus-eq, simp add:aGroup.ag-pOp-commute[of K -a x],
    frule val-nonzero-z[of v x], assumption+, erule exE,
    simp add:a-zpz aminus, simp add:ant-0[THEN sym] aless-zless,
    assumption)
done

lemma (in Corps) Ostrowski-first:vals-nonequiv K (Suc 0) vv
  ⇒ ∃ x ∈ carrier K. Ostrowski-elem K (Suc 0) vv x
  apply (simp add:vals-nonequiv-def,
    cut-tac Nset-Suc0, (erule conjE)+,
    simp add:valuations-def)
  apply (rotate-tac -1,

```

```

    frule-tac a = 0 in forall-spec, simp,
    rotate-tac -1,
    drule-tac a = Suc 0 in forall-spec, simp)
  apply (drule-tac a = Suc 0 in forall-spec, simp,
    rotate-tac -1,
    drule-tac a = 0 in forall-spec, simp, simp)
  apply (frule-tac a = 0 in forall-spec, simp,
    drule-tac a = Suc 0 in forall-spec, simp,
    frule-tac v = vv 0 and v' = vv (Suc 0) in
      nonequiv-ex-Ostrowski-elem, assumption+,
    erule bexE)

  apply (erule conjE,
    frule-tac v = vv (Suc 0) and v' = vv 0 in
      nonequiv-ex-Ostrowski-elem, assumption+,
    erule bexE,
    thin-tac  $\neg$  v-equiv K (vv (Suc 0)) (vv 0),
    thin-tac  $\neg$  v-equiv K (vv 0) (vv (Suc 0)))

  apply (rename-tac s t)
  apply (erule conjE,
    frule-tac x = t and v = vv 0 in val-neg-nonzero, assumption+)
  apply (simp add:less-ant-def, (erule conjE)+,
    frule-tac x = s and v = vv (Suc 0) in val-neg-nonzero,
    assumption+, unfold less-ant-def)
  apply (rule conjI, assumption+)

  apply (frule-tac s = s and t = t and v = vv 0 in OstrowskiTr2,
    assumption+, rule ale-neq-less, assumption+)
  apply (frule-tac s = s and t = t and v = vv (Suc 0) in OstrowskiTr3,
    assumption+, rule ale-neq-less, assumption+)
  apply (subgoal-tac t  $\cdot_r$  ((s  $\pm$  t)-K)  $\in$  carrier K,
    simp only:Ostrowski-elem-def,
    simp only: nset-m-m[of Suc 0], blast)

  apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    rule Ring.ring-tOp-closed, assumption+,
    frule-tac s = s and t = t and v = vv 0 in OstrowskiTr1,
    assumption+, rule ale-neq-less, assumption+,
    frule-tac x = s and y = t in aGroup.ag-pOp-closed[of K], assumption+)
  apply (cut-tac x = s  $\pm$  t in invf-closed, blast)
  apply assumption
done

```

```

lemma (in Corps) Ostrowski: $\forall$  vv. vals-nonequiv K (Suc n) vv  $\longrightarrow$ 
  ( $\exists x \in \text{carrier } K. \text{Ostrowski-elem } K (Suc n) vv x$ )
  apply (induct-tac n,

```

rule allI, rule impI, simp add:Ostrowski-first)

apply (*rule allI, rule impI,*
frule-tac n = n and vv = vv in restrict-vals-nonequiv1,
frule-tac n = n and vv = vv in restrict-vals-nonequiv2,
frule-tac a = compose {h. h ≤ Suc n} vv (skip 1) in forall-spec,
assumption,
drule-tac a = compose {h. h ≤ Suc n} vv (skip 2) in forall-spec,
assumption+, (erule bexE)+)
apply (*rename-tac n vv s t,*
cut-tac field-is-ring, frule Ring.ring-is-ag[of K])

apply (*frule-tac x = s and y = t in Ring.ring-tOp-closed[of K], assumption+,*
case-tac 0 ≤ vv (Suc 0) s ∧ 0 ≤ vv (Suc (Suc 0)) t,
frule-tac vv = vv and s = s and t = t in OstrowskiTr5, assumption+)
apply *blast*

apply (*simp,*
case-tac 0 ≤ (vv (Suc 0) s), simp,
frule-tac n = Suc (Suc n) and m = Suc (Suc 0) and vv = vv in
vals-nonequiv-valuation,
simp,
frule-tac v = vv (Suc (Suc 0)) and x = t in OstrowskiTr6,
assumption+,
frule-tac x = 1_r ± t ·_r (1_r ± -_a t) in invf-closed1,
frule-tac x = t and y = (1_r ± t ·_r (1_r ± -_a t))^{-K} in
Ring.ring-tOp-closed, assumption+, simp)
apply (*subgoal-tac Ostrowski-elem K (Suc (Suc n)) vv*
(t ·_r ((1_r ± t ·_r (1_r ± -_a t))^{-K})),
blast)
apply (*subst Ostrowski-elem-def,*
rule conjI,
thin-tac Ostrowski-elem K (Suc n)
(compose {h. h ≤ Suc n} vv (skip (Suc 0))) s,
thin-tac vals-nonequiv K (Suc n)
(compose {h. h ≤ Suc n} vv (skip (Suc 0))),
thin-tac vals-nonequiv K (Suc n) (compose {h. h ≤ Suc n} vv (skip 2)),
thin-tac 0 ≤ (vv (Suc 0) s),
frule-tac n = Suc (Suc n) and vv = vv and m = 0 in
vals-nonequiv-valuation, simp,
rule-tac v = vv 0 and x = t in
OstrowskiTr8, assumption+)
apply (*simp add:Ostrowski-elem-def, (erule conjE)+,*
thin-tac ∀j∈nset (Suc 0) (Suc n).
0 < (compose {h. h ≤ (Suc n)} vv (skip 2) j t),
simp add:compose-def skip-def,

$rule\ ballI,$
 $thin-tac\ 0 \leq (vv\ (Suc\ 0)\ s),$
 $thin-tac\ Ostrowski-elem\ K\ (Suc\ n)$
 $\quad (compose\ \{h.\ h \leq (Suc\ n)\}\ vv\ (skip\ (Suc\ 0)))\ s,$
 $frule-tac\ n = Suc\ (Suc\ n)\ \text{and}\ vv = vv\ \text{and}\ m = j\ \text{in}$
 $\quad vals-nonequiv-valuation,$
 $simp\ add:nset-def,\ simp\ add:Ostrowski-elem-def,\ (erule\ conjE)+)$

$apply\ (case-tac\ j = Suc\ 0,\ simp,$
 $\quad drule-tac\ x = Suc\ 0\ \text{in}\ bspec,$
 $\quad simp\ add:nset-def,$
 $\quad simp\ add:compose-def\ skip-def,$
 $\quad rule-tac\ v = vv\ (Suc\ 0)\ \text{and}\ x = t\ \text{in}$
 $\quad\quad OstrowskiTr9,\ assumption+,$
 $\quad frule-tac\ j = j\ \text{and}\ n = n\ \text{in}\ nset-Tr51,\ assumption+,$
 $\quad drule-tac\ x = j - Suc\ 0\ \text{in}\ bspec,\ assumption+,$
 $\quad simp\ add:compose-def\ skip-def)$

$apply\ (case-tac\ j = Suc\ (Suc\ 0),\ simp)\ apply\ ($
 $\quad rule-tac\ v = vv\ (Suc\ (Suc\ 0))\ \text{and}\ x = t\ \text{in}\ OstrowskiTr10,$
 $\quad assumption+)\ apply\ ($
 $\quad subgoal-tac\ \neg j - Suc\ 0 \leq Suc\ 0,\ simp\ add:nset-def)\ apply\ ($
 $\quad rule-tac\ v = vv\ j\ \text{and}\ x = t\ \text{in}$
 $\quad\quad OstrowskiTr9)\ apply\ (simp\ add:nset-def,\ assumption+)$

$apply\ (simp\ add:nset-def,\ (erule\ conjE)+,\ rule\ nset-Tr52,\ assumption+,$
 $\quad thin-tac\ vals-nonequiv\ K\ (Suc\ n)$
 $\quad\quad (compose\ \{h.\ h \leq (Suc\ n)\}\ vv\ (skip\ (Suc\ 0))),$
 $\quad thin-tac\ vals-nonequiv\ K\ (Suc\ n)$
 $\quad\quad (compose\ \{h.\ h \leq (Suc\ n)\}\ vv\ (skip\ 2)),$
 $\quad thin-tac\ Ostrowski-elem\ K\ (Suc\ n)$
 $\quad\quad (compose\ \{h.\ h \leq (Suc\ n)\}\ vv\ (skip\ 2))\ t)$

$apply\ (subgoal-tac\ s \cdot_r ((1_r \pm s \cdot_r (1_r \pm -_a s))^{-K}) \in carrier\ K,$
 $\quad subgoal-tac\ Ostrowski-elem\ K\ (Suc\ (Suc\ n))\ vv$
 $\quad\quad (s \cdot_r ((1_r \pm s \cdot_r (1_r \pm -_a s))^{-K})),\ blast)$

$prefer\ 2$
 $apply\ (frule-tac\ n = Suc\ (Suc\ n)\ \text{and}\ m = Suc\ 0\ \text{and}\ vv = vv\ \text{in}$
 $\quad vals-nonequiv-valuation,\ simp,$
 $\quad frule-tac\ v = vv\ (Suc\ 0)\ \text{and}\ x = s\ \text{in}\ OstrowskiTr6,\ assumption+,$
 $\quad rule\ Ring.ring-tOp-closed,\ assumption+,$
 $\quad frule-tac\ x = 1_r \pm s \cdot_r (1_r \pm -_a s)\ \text{in}\ invf-closed1,\ simp,$
 $\quad simp\ add:Ostrowski-elem-def)$

$apply\ (rule\ conjI)$
 $apply\ (rule-tac\ v = vv\ 0\ \text{and}\ x = s\ \text{in}\ OstrowskiTr8,$
 $\quad simp\ add:vals-nonequiv-valuation,\ assumption+)\ apply\ ($
 $\quad thin-tac\ vals-nonequiv\ K\ (Suc\ (Suc\ n))\ vv,$
 $\quad (erule\ conjE)+,$
 $\quad thin-tac\ \forall j \in nset\ (Suc\ 0)\ (Suc\ n).$
 $\quad\quad 0 < (compose\ \{h.\ h \leq (Suc\ n)\}\ vv\ (skip\ (Suc\ 0))\ j\ s),$

$\text{simp add:compose-def skip-def, rule ballI})$

apply (*case-tac* $j = \text{Suc } 0$, *simp*,
rule-tac $v = vv (\text{Suc } 0)$ **and** $x = s$ **in** *OstrowskiTr10*,
simp add:vals-nonequiv-valuation, assumption+,
rule-tac $v = vv j$ **and** $x = s$ **in** *OstrowskiTr9*,
simp add:vals-nonequiv-valuation nset-def, assumption,
(erule conjE)+, simp add:compose-def skip-def,
frule-tac $j = j$ **in** *nset-Tr51, assumption+*,
drule-tac $x = j - \text{Suc } 0$ **in** *bspec, assumption+*)
apply (*simp add:nset-def*)
done

lemma (**in** *Corps*) *val-1-nonzero*: $\llbracket \text{valuation } K \ v; x \in \text{carrier } K; v \ x = 1 \rrbracket \implies$
 $x \neq \mathbf{0}$

apply (*rule contrapos-pp, simp+*,
simp add:value-of-zero,
rotate-tac 3, drule sym, simp only:ant-1[THEN sym],
simp del:ant-1)
done

lemma (**in** *Corps*) *Approximation1-5Tr1*: $\llbracket \text{vals-nonequiv } K (\text{Suc } n) \ vv;$
 $n\text{-val } K (vv \ 0) = vv \ 0; a \in \text{carrier } K; vv \ 0 \ a = 1; x \in \text{carrier } K;$
 $\text{Ostrowski-elem } K (\text{Suc } n) \ vv \ x \rrbracket \implies$
 $\forall m. 2 \leq m \longrightarrow vv \ 0 \ ((1_r \pm -_a \ x)^{\wedge K \ m} \pm a \cdot_r (x^{\wedge K \ m})) = 1$

apply (*cut-tac field-is-ring, frule Ring.ring-is-ag[of K],*
frule Ring.ring-one[of K],
rule allI, rule impI,
frule vals-nonequiv-valuation[of Suc n vv 0],
simp,
simp add:Ostrowski-elem-def, frule conjunct1, fold Ostrowski-elem-def,
frule val-1-nonzero[of vv 0 a], assumption+)
apply (*frule vals-nonequiv-valuation[of Suc n vv 0], simp,*
frule val-nonzero-noninf[of vv 0 a], assumption+,
frule val-unit-cond[of vv 0 x], assumption+,
frule-tac $n = m$ **in** *Ring.npClose[of K x], assumption+,*
frule aGroup.ag-mOp-closed[of K x], assumption+,
frule aGroup.ag-pOp-closed[of K 1_r -_a x], assumption+)
apply (*subgoal-tac* $0 < m$,
frule-tac $x = a \cdot_r (x^{\wedge K \ m})$ **and** $y = (1_r \pm -_a \ x)^{\wedge K \ m}$ **in**
value-less-eq[of vv 0],
rule Ring.ring-tOp-closed, assumption+,
rule Ring.npClose, assumption+, simp add: val-t2p,
frule value-zero-nonzero[of vv 0 x], assumption+,
simp add:val-exp-ring[THEN sym], simp add:asprod-n-0 aadd-0-r)
apply (*case-tac* $x = 1_r K$, *simp add:aGroup.ag-r-inv1,*
frule-tac $n = m$ **in** *Ring.npZero-sub[of K], simp,*
simp add:value-of-zero)
apply (*cut-tac inf-ge-any[of 1], simp add: less-le*)

apply (*rotate-tac* -1 , *drule* *not-sym*,
frule *aGroup.ag-neq-diffnonzero*[*of K 1_r x*],
simp *add:Ring.ring-one*[*of K*], *assumption+*, *simp*,
simp *add:val-exp-ring*[*THEN sym*],
cut-tac $n1 = m$ **in** *of-nat-0-less-iff*[*THEN sym*])
apply (*cut-tac* $a = 0 < m$ **and** $b = 0 < \text{int } m$ **in** *a-b-exchange*, *simp*,
assumption)
apply (*thin-tac* $(0 < m) = (0 < \text{int } m)$,
frule *val-nonzero-z*[*of vv 0 1_r $\pm -_a x$*], *assumption+*,
erule *exE*, *simp*, *simp* *add:asprod-amult* $a-z-z$,
simp *only:ant-1*[*THEN sym*], *simp* *only:aless-zless*, *simp* *add:ge2-zmult-pos*)

apply (*subst* *aGroup.ag-pOp-commute*[*of K*], *assumption+*,
rule *Ring.npClose*, *assumption+*, *rule* *Ring.ring-tOp-closed*[*of K*],
assumption+,
rotate-tac -1 , *drule* *sym*, *simp*,
thin-tac $vv\ 0\ (a \cdot_r x^{\wedge K m} \pm (1_r \pm -_a x)^{\wedge K m}) = vv\ 0\ (a \cdot_r x^{\wedge K m})$)
apply (*simp* *add:val-t2p*,
frule *value-zero-nonzero*[*of vv 0 x*], *assumption+*,
simp *add:val-exp-ring*[*THEN sym*], *simp* *add:asprod-n-0*,
simp *add:aadd-0-r*,
cut-tac $z = m$ **in** *less-le-trans*[*of 0 2*], *simp*, *assumption+*)
done

lemma (**in** *Corps*) *Approximation1-5Tr3*: $\llbracket \text{vals-nonequiv } K\ (\text{Suc } n)\ vv;$
 $x \in \text{carrier } K; \text{Ostrowski-elem } K\ (\text{Suc } n)\ vv\ x; j \in \text{nset } (\text{Suc } 0)\ (\text{Suc } n) \rrbracket$
 $\implies vv\ j\ ((1_r \pm -_a x)^{\wedge K m}) = 0$
apply (*frule* *Ostrowski-elem-not-one*[*of n vv x*], *assumption+*,
cut-tac *field-is-ring*, *frule* *Ring.ring-is-ag*[*of K*],
frule *Ring.ring-one*[*of K*],
frule *aGroup.ag-pOp-closed*[*of K 1_r -_a x*], *assumption+*)
apply (*simp* *add:aGroup.ag-mOp-closed*, *simp* *add:nset-def*,
frule-tac $m = j$ **in** *vals-nonequiv-valuation*[*of Suc n vv*],
simp,
frule-tac $v1 = vv\ j$ **and** $x1 = 1_r \pm -_a x$ **and** $n1 = m$ **in**
val-exp-ring[*THEN sym*], *assumption+*)

apply (*frule-tac* $v = vv\ j$ **and** $x = 1_r$ **and** $y = -_a x$ **in**
value-less-eq, *assumption+*, *simp* *add:aGroup.ag-mOp-closed*)
apply (*simp* *add:value-of-one*, *simp* *add:val-minus-eq*,
simp *add:Ostrowski-elem-def* *nset-def*)
apply (*simp* *add:value-of-one*, *rotate-tac* -1 , *drule* *sym*,
simp *add:asprod-n-0*)
done

lemma (**in** *Corps*) *Approximation1-5Tr4*: $\llbracket \text{vals-nonequiv } K\ (\text{Suc } n)\ vv;$
 $aa \in \text{carrier } K; x \in \text{carrier } K;$
 $\text{Ostrowski-elem } K\ (\text{Suc } n)\ vv\ x; j \leq (\text{Suc } n) \rrbracket \implies$
 $vv\ j\ (aa \cdot_r (x^{\wedge K m})) = vv\ j\ aa + (\text{int } m) *_a (vv\ j\ x)$

$value-less-eq[of\ vv\ j],$
 $rule\ Ring.npClose,\ assumption+,$
 $rule\ aGroup.ag-pOp-closed,\ assumption+, simp\ add:Ring.ring-one,$
 $simp\ add:aGroup.ag-mOp-closed)$
apply ($rule\ Ring.ring-tOp-closed,\ assumption+,$
 $rule\ Ring.npClose,\ assumption+,$
 $simp\ add:Approximation1-5Tr3,$
 $frule\ sym,\ assumption)$
done

lemma (**in** *Corps*) *Approximation1-5Tr7*: $\llbracket a \in carrier\ K; vv\ 0\ a = 1;$
 $x \in carrier\ K \rrbracket \implies$
 $vals-nonequiv\ K\ (Suc\ n)\ vv \wedge Ostrowski-elem\ K\ (Suc\ n)\ vv\ x \longrightarrow$
 $(\exists\ l.\ \forall\ m.\ l < m \longrightarrow (\forall\ j \in nset\ (Suc\ 0)\ (Suc\ n). \$
 $(vv\ j\ ((1_r \pm -_a\ x) \frown^K m \pm a \cdot_r (x \frown^K m)) = 0)))$

apply ($induct-tac\ n,$
 $rule\ impI,\ erule\ conjE,\ simp\ add:nset-m-m[of\ Suc\ 0],$
 $frule\ vals-nonequiv-valuation[of\ Suc\ 0\ vv\ Suc\ 0],\ simp,$
 $frule\ Approximation1-5Tr6[of\ 0\ vv\ a\ x\ Suc\ 0],\ assumption+)$
apply ($frule\ vals-nonequiv-valuation[of\ Suc\ 0\ vv\ 0],\ simp,$
 $frule\ val-1-nonzero[of\ vv\ 0\ a],\ assumption+, simp\ add:nset-def,$
 $assumption)$

apply ($rule\ impI,\ erule\ conjE,$
 $frule-tac\ n = n\ \mathbf{in}\ restrict-vals-nonequiv[of\ -\ vv],$
 $frule-tac\ n = n\ \mathbf{in}\ restrict-Ostrowski-elem[of\ x\ -\ vv],$
 $assumption,\ simp,$
 $erule\ exE,$
 $frule-tac\ n = Suc\ n\ \mathbf{and}\ j = Suc\ (Suc\ n)\ \mathbf{in}\ Approximation1-5Tr6$
 $[of\ -\ vv\ a\ x],\ assumption+,$
 $frule-tac\ n = Suc\ (Suc\ n)\ \mathbf{in}\ vals-nonequiv-valuation[of\ -\ vv$
 $0],\ simp,$
 $rule\ val-1-nonzero[of\ vv\ 0\ a],\ assumption+,$
 $simp\ add:nset-def)$

apply ($erule\ exE,$
 $subgoal-tac\ \forall\ m.\ (max\ l\ la) < m \longrightarrow (\forall\ j \in nset\ (Suc\ 0)\ (Suc\ (Suc\ n)).$
 $vv\ j\ ((1_r \pm -_a\ x) \frown^K m \pm a \cdot_r (x \frown^K m)) = 0),$
 $blast,$
 $simp\ add:nset-Suc)$
done

lemma (**in** *Corps*) *Approximation1-5P*: $\llbracket vals-nonequiv\ K\ (Suc\ n)\ vv;$
 $n-val\ K\ (vv\ 0) = vv\ 0 \rrbracket \implies$
 $\exists\ x \in carrier\ K.\ ((vv\ 0\ x = 1) \wedge (\forall\ j \in nset\ (Suc\ 0)\ (Suc\ n). (vv\ j\ x) = 0))$
apply ($frule\ vals-nonequiv-valuation[of\ Suc\ n\ vv\ 0],\ simp)$ **apply** ($frule\ n-val-surj[of\ vv\ 0],\ erule\ bexE)$ **apply** ($rename-tac\ aa)$ **apply** ($cut-tac\ n = n\ \mathbf{in}\ Ostrowski)$ **apply** ($drule-tac\ a = vv\ \mathbf{in}\ forall-spec[of\ vals-nonequiv\ K\ (Suc\ n)],\ simp)$

```

apply (
  erule bexE,
  frule-tac a = aa and x = x in Approximation1-5Tr1[of n vv],
  assumption+,
  simp, assumption+)
apply (frule-tac a = aa and x = x in Approximation1-5Tr7[of - vv - n],
  simp, assumption,
  simp, erule exE,
  cut-tac b = Suc l in max.cobounded1[of 2],
  cut-tac b = Suc l in max.cobounded2[of - 2],
  cut-tac n = l in lessI,
  frule-tac x = l and y = Suc l and z = max 2 (Suc l) in
    less-le-trans, assumption+,
  thin-tac Suc l ≤ max 2 (Suc l), thin-tac l < Suc l,
  drule-tac a = max 2 (Suc l) in forall-spec, simp,
  drule-tac a = max 2 (Suc l) in forall-spec, assumption)
apply (subgoal-tac (1r ± -a x)~K (max 2 (Suc l)) ± aa ·r (x~K (max 2 (Suc l))))
  ∈
    carrier K,
    blast,
    cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    rule aGroup.ag-pOp-closed, assumption+, rule Ring.npClose, assumption+,
    rule aGroup.ag-pOp-closed, assumption+, simp add:Ring.ring-one,
    simp add:aGroup.ag-mOp-closed, rule Ring.ring-tOp-closed, assumption+,
    rule Ring.npClose, assumption+)
done

lemma K-gamma-hom:k ≤ n ⇒ ∀ j ≤ n. (λl. γk l) j ∈ Zset
apply (simp add:Zset-def)
done

lemma transpos-eq:(τ0 0) k = k
by (simp add:transpos-def)

lemma (in Corps) transpos-vals-nonequiv:⌈vals-nonequiv K (Suc n) vv;
  j ≤ (Suc n)⌋ ⇒ vals-nonequiv K (Suc n) (vv ∘ (τ0 j))
apply (simp add:vals-nonequiv-def)
apply (frule conjunct1, fold vals-nonequiv-def)
apply (simp add:valuations-def, rule conjI)
apply (rule allI, rule impI)
apply (case-tac ja = 0, simp,
  case-tac j = 0, simp add:transpos-eq)
apply (subst transpos-ij-1[of 0 Suc n j], simp, assumption+,
  rule not-sym, assumption, simp)

apply (case-tac ja = j, simp)
apply (subst transpos-ij-2[of 0 Suc n j], simp, assumption, simp,
  simp add:vals-nonequiv-valuation)
apply (case-tac j = 0, simp add:transpos-eq)

```

```

apply (cut-tac  $x = ja$  in transpos-id-1[of 0 Suc n j], simp, assumption+,
      rule not-sym, assumption+)
apply (simp add:vals-nonequiv-valuation,
      (rule allI, rule impI)+, rule impI)
apply (case-tac  $j = 0$ , simp add:transpos-eq,
      simp add:vals-nonequiv-def,
      cut-tac transpos-inj[of 0 Suc n j], simp)
apply (frule-tac  $x = ja$  and  $y = l$  in injective[of transpos 0 j
      { $j. j \leq (Suc\ n)$ }], simp, simp, assumption+)
apply (cut-tac  $l = ja$  in transpos-mem[of 0 Suc n j], simp, assumption+,
      simp, assumption,
      cut-tac  $l = l$  in transpos-mem[of 0 Suc n j], simp, assumption+,
      simp, assumption)
apply (simp add:vals-nonequiv-def,
      simp, assumption, rule not-sym, assumption)
done

```

definition

$Ostrowski-base :: [-, nat \Rightarrow 'b \Rightarrow ant, nat] \Rightarrow (nat \Rightarrow 'b)$
 $(\langle (\Omega_- _ -) \rangle [90,90,91]90) \textbf{ where}$
 $Ostrowski-base\ K\ vv\ n = (\lambda j \in \{h. h \leq n\}. (SOME\ x. x \in carrier\ K \wedge$
 $(Ostrowski-elem\ K\ n\ (vv \circ (\tau_0\ j))\ x)))$

definition

$App-base :: [-, nat \Rightarrow 'b \Rightarrow ant, nat] \Rightarrow (nat \Rightarrow 'b) \textbf{ where}$
 $App-base\ K\ vv\ n = (\lambda j \in \{h. h \leq n\}. (SOME\ x. x \in carrier\ K \wedge (((vv \circ \tau_0\ j)\ 0\ x$
 $= 1) \wedge (\forall k \in nset\ (Suc\ 0)\ n. ((vv \circ \tau_0\ j)\ k\ x) = 0))))$

lemma (**in** Corps) $Ostrowski-base-hom:vals-nonequiv\ K\ (Suc\ n)\ vv \Longrightarrow$

$Ostrowski-base\ K\ vv\ (Suc\ n) \in \{h. h \leq (Suc\ n)\} \rightarrow carrier\ K$

```

apply (rule Pi-I, rename-tac l,
      simp add:Ostrowski-base-def,
      frule-tac  $j = l$  in transpos-vals-nonequiv[of n vv], simp,
      cut-tac Ostrowski[of n])

```

```

apply (drule-tac  $a = vv \circ \tau_0\ l$  in forall-spec, simp,
      rule someI2-ex, blast, simp)

```

done

lemma (**in** Corps) $Ostrowski-base-mem:vals-nonequiv\ K\ (Suc\ n)\ vv \Longrightarrow$

$\forall j \leq (Suc\ n). Ostrowski-base\ K\ vv\ (Suc\ n)\ j \in carrier\ K$

```

by (rule allI, rule impI,
      frule Ostrowski-base-hom[of n vv],
      simp add:funcset-mem del:Pi-I')

```

lemma (**in** Corps) $Ostrowski-base-mem-1: \llbracket vals-nonequiv\ K\ (Suc\ n)\ vv;$

$j \leq (Suc\ n) \rrbracket \Longrightarrow Ostrowski-base\ K\ vv\ (Suc\ n)\ j \in carrier\ K$

by (simp add:Ostrowski-base-mem)

lemma (in Corps) Ostrowski-base-nonzero: $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv; j \leq \text{Suc } n \rrbracket \implies (\Omega_K \text{ } vv \text{ (Suc } n)) \text{ } j \neq 0$

apply (simp add: Ostrowski-base-def,
frule-tac $j = j$ in transpos-vals-nonequiv[of n vv],
assumption+,
cut-tac Ostrowski[of n],
drule-tac $a = vv \circ \tau_0 j$ in forall-spec, assumption,
rule someI2-ex, blast)

apply (thin-tac $\exists x \in \text{carrier } K. \text{ Ostrowski-elem } K \text{ (Suc } n) (vv \circ \tau_0 j) \text{ } x,$
erule conjE)

apply (rule-tac $vv = vv \circ \tau_0 j$ and $x = x$ in Ostrowski-elem-nonzero[of n],
assumption+)

done

lemma (in Corps) Ostrowski-base-pos: $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv; j \leq \text{Suc } n; ja \leq \text{Suc } n; ja \neq j \rrbracket \implies 0 < ((vv \text{ } j) ((\Omega_K \text{ } vv \text{ (Suc } n)) \text{ } ja))$

apply (simp add: Ostrowski-base-def,
frule-tac $j = ja$ in transpos-vals-nonequiv[of n vv],
assumption+,
cut-tac Ostrowski[of n],
drule-tac $a = vv \circ \tau_0 ja$ in forall-spec, assumption+)

apply (rule someI2-ex, blast,
thin-tac $\exists x \in \text{carrier } K. \text{ Ostrowski-elem } K \text{ (Suc } n) (vv \circ \tau_0 ja) \text{ } x,$
simp add: Ostrowski-elem-def, (erule conjE)+)

apply (case-tac $ja = 0$, simp, cut-tac transpos-eq[of j],
simp add: nset-def, frule Suc-leI[of $0 \text{ } j$],
frule-tac $a = j$ in forall-spec, simp, simp)

apply (case-tac $j = 0$, simp,
frule-tac $x = ja$ in bspec, simp add: nset-def,
cut-tac transpos-ij-2[of $0 \text{ Suc } n \text{ } ja$], simp, simp+)

apply (frule-tac $x = j$ in bspec, simp add: nset-def,
cut-tac transpos-id[of $0 \text{ Suc } n \text{ } ja \text{ } j$], simp+)

done

lemma (in Corps) App-base-hom: $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv; \forall j \leq (\text{Suc } n). \text{ n-val } K \text{ (} vv \text{ } j) = vv \text{ } j \rrbracket \implies \forall j \leq (\text{Suc } n). \text{ App-base } K \text{ } vv \text{ (Suc } n) \text{ } j \in \text{carrier } K$

apply (rule allI, rule impI,
rename-tac k ,
subst App-base-def)

apply (case-tac $k = 0$, simp, simp add: transpos-eq,
frule Approximation1-5P[of n vv], simp,
rule someI2-ex, blast, simp)

apply (frule-tac $j = k$ in transpos-vals-nonequiv[of n vv],
simp add: nset-def,
frule-tac $vv = vv \circ \tau_0 k$ in Approximation1-5P[of n])

apply (simp add: cmp-def, subst transpos-ij-1[of $0 \text{ Suc } n$], simp+,
subst transpos-ij-1[of $0 \text{ Suc } n \text{ } k$], simp+)

apply (rule someI2-ex, blast, simp)

done

lemma (in Corps) Approximation1-5P2: $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv; \forall l \in \{h. h \leq \text{Suc } n\}. n\text{-val } K \text{ (} vv \text{ } l) = vv \text{ } l; i \leq \text{Suc } n; j \leq \text{Suc } n \rrbracket \implies vv \text{ } i \text{ (App-base } K \text{ } vv \text{ (Suc } n) \text{ } j) = \delta_i j$

apply (simp add: App-base-def)

apply (case-tac $j = 0$, simp add: transpos-eq, rule someI2-ex, frule Approximation1-5P[of $n \text{ } vv$], simp, blast, simp add: Kronecker-delta-def, rule impI, (erule conjE)+, frule-tac $x = i$ in bspec, simp add: nset-def, assumption)

apply (frule-tac $j = j$ in transpos-vals-nonequiv[of $n \text{ } vv$], simp, frule Approximation1-5P[of $n \text{ } vv \circ \tau_0 j$], simp add: cmp-def, simp add: transpos-ij-1[of $0 \text{ Suc } n \text{ } j$])

apply (simp add: cmp-def, case-tac $i = 0$, simp add: transpos-eq, simp add: transpos-ij-1, simp add: Kronecker-delta-def, rule someI2-ex, blast, thin-tac $\exists x \in \text{carrier } K. vv \text{ } j \text{ } x = 1 \wedge (\forall ja \in \text{nset (Suc } 0) \text{ (Suc } n). vv \text{ ((}\tau_0 j) \text{ } ja) \text{ } x = 0)$, (erule conjE)+, drule-tac $x = j$ in bspec, simp add: nset-def, simp add: transpos-ij-2)

apply (simp add: Kronecker-delta-def, case-tac $i = j$, simp add: transpos-ij-1, rule someI2-ex, blast, simp)

apply (simp, rule someI2-ex, blast, thin-tac $\exists x \in \text{carrier } K. vv \text{ ((}\tau_0 j) \text{ } 0) \text{ } x = 1 \wedge (\forall ja \in \text{nset (Suc } 0) \text{ (Suc } n). vv \text{ ((}\tau_0 j) \text{ } ja) \text{ } x = 0)$, (erule conjE)+, drule-tac $x = i$ in bspec, simp add: nset-def, cut-tac transpos-id[of $0 \text{ Suc } n \text{ } j \text{ } i$], simp+)

done

lemma (in Corps) Approximation1-5: $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv; \forall j \leq \text{Suc } n. n\text{-val } K \text{ (} vv \text{ } j) = vv \text{ } j \rrbracket \implies \exists x. (\forall j \leq \text{Suc } n. x \text{ } j \in \text{carrier } K) \wedge (\forall i \leq \text{Suc } n. \forall j \leq \text{Suc } n. ((vv \text{ } i) \text{ (} x \text{ } j) = \delta_i j))$

apply (frule App-base-hom[of $n \text{ } vv$], rule allI, simp)

apply (subgoal-tac $(\forall i \leq \text{Suc } n. \forall j \leq \text{Suc } n. (vv \text{ } i) \text{ ((App-base } K \text{ } vv \text{ (Suc } n) \text{ } j) = (\delta_i j)))$, blast)

apply (rule allI, rule impI)+

apply (rule Approximation1-5P2, assumption+, simp+)

done

lemma (in Corps) Ostrowski-baseTr0: $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv; l \leq (\text{Suc } n) \rrbracket$
 $\implies 0 < ((vv \ l) \ (1_r \pm -_a \ (Ostrowski-base \ K \ vv \ (\text{Suc } n) \ l))) \wedge$
 $(\forall m \in \{h. h \leq (\text{Suc } n)\} - \{l\}. \ 0 < ((vv \ m) \ (Ostrowski-base \ K \ vv \ (\text{Suc } n) \ l)))$

apply (simp add: Ostrowski-base-def,
 frule-tac $j = l$ in transpos-vals-nonequiv[of $n \ vv$], assumption,
 cut-tac Ostrowski[of n],
 drule-tac $a = vv \circ \tau_0 \ l$ in forall-spec, assumption)

apply (erule bexE,
 unfold Ostrowski-elem-def, frule conjunct1,
 fold Ostrowski-elem-def,
 rule conjI, simp add: Ostrowski-elem-def)

apply (case-tac $l = 0$, simp, simp add: transpos-eq,
 rule someI2-ex, blast, simp,
 simp add: transpos-ij-1,
 rule someI2-ex, blast, simp)

apply (simp add: Ostrowski-elem-def,
 case-tac $l = 0$, simp, simp add: transpos-eq,
 rule someI2-ex, blast,
 thin-tac $0 < vv \ 0 \ (1_r \pm -_a \ x) \wedge$
 $(\forall j \in \text{nset } (\text{Suc } 0) \ (\text{Suc } n). \ 0 < vv \ j \ x),$
 rule ballI, simp add: nset-def)

apply (rule ballI, erule conjE,
 rule someI2-ex, blast,
 thin-tac $\forall j \in \text{nset } (\text{Suc } 0) \ (\text{Suc } n). \ 0 < vv \ ((\tau_0 \ l) \ j) \ x,$
 (erule conjE)+)

apply (case-tac $m = 0$, simp,
 drule-tac $x = l$ in bspec, simp add: nset-def,
 simp add: transpos-ij-2,
 drule-tac $x = m$ in bspec, simp add: nset-def,
 simp add: transpos-id)

done

lemma (in Corps) Ostrowski-baseTr1: $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv; l \leq (\text{Suc } n) \rrbracket$
 $\implies 0 < ((vv \ l) \ (1_r \pm -_a \ (Ostrowski-base \ K \ vv \ (\text{Suc } n) \ l)))$

by (simp add: Ostrowski-baseTr0)

lemma (in Corps) Ostrowski-baseTr2: $\llbracket \text{vals-nonequiv } K \text{ (Suc } n) \text{ } vv;$
 $l \leq (\text{Suc } n); m \leq (\text{Suc } n); l \neq m \rrbracket \implies$
 $0 < ((vv \ m) \ (Ostrowski-base \ K \ vv \ (\text{Suc } n) \ l))$

apply (frule Ostrowski-baseTr0[of $n \ vv \ l$], assumption+)

apply simp

done

lemma Nset-have-two: $j \in \{h. h \leq (\text{Suc } n)\} \implies \exists m \in \{h. h \leq (\text{Suc } n)\}. j \neq m$

```

apply (rule contrapos-pp, simp+,
        case-tac j = Suc n, simp,
        drule-tac a = 0 in forall-spec, simp, arith)
apply (drule-tac a = Suc n in forall-spec, simp, simp)
done

lemma (in Corps) Ostrowski-base-mpow-not-one:  $\llbracket 0 < N; j \leq \text{Suc } n; \text{vals-nonequiv } K (\text{Suc } n) \text{ vv} \rrbracket \implies$ 

$$1_r \pm -_a ((\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N) \neq \mathbf{0}$$

apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
        rule contrapos-pp, simp+,
        frule Ostrowski-base-mem-1[of n vv j], assumption,
        frule Ring.npClose[of K  $(\Omega_K \text{ vv } (\text{Suc } n)) j N$ ], assumption+,
        frule Ring.ring-one[of K],
        frule aGroup.ag-mOp-closed[of K  $(\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N$ ], assumption+,
        frule aGroup.ag-pOp-closed[of K  $1_r -_a ((\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N)$ ],
        assumption+)
apply (frule aGroup.ag-pOp-add-r[of K  $1_r \pm -_a ((\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N) \mathbf{0}$ 
         $(\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N$ ], assumption+,
        simp add:aGroup.ag-inc-zero, assumption+,
        thin-tac  $1_r \pm -_a ((\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N) = \mathbf{0}$ )
apply (simp add:aGroup.ag-pOp-assoc[of K  $1_r -_a ((\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N)$ ])
apply (simp add:aGroup.ag-l-inv1, simp add:aGroup.ag-r-zero aGroup.ag-l-zero)
apply (subgoal-tac  $\forall m \leq (\text{Suc } n). (j \neq m \longrightarrow$ 
 $0 < (\text{vv } m ((\Omega_K \text{ vv } (\text{Suc } n)) j))))$ 
apply (cut-tac Nset-have-two[of j n],
        erule bexE, drule-tac a = m in forall-spec, simp,
        thin-tac  $(\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N \pm -_a ((\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N) \in \text{carrier } K,$ 
        frule-tac f = vv m in eq-elems-eq-val[of  $1_r (\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N$ ],
        thin-tac  $1_r = (\Omega_K \text{ vv } (\text{Suc } n)) j^{\sim K} N$ , simp)
apply (frule-tac m = m in vals-nonequiv-valuation[of Suc n vv],
        assumption+,
        frule-tac v1 = vv m and n1 = N in val-exp-ring[THEN sym,
        of  $- (\Omega_K \text{ vv } (\text{Suc } n)) j$ ], assumption+,
        simp add:Ostrowski-base-nonzero, simp, simp add:value-of-one)
apply (subgoal-tac int N  $\neq 0$ ,
        frule-tac x = vv m  $((\Omega_K \text{ vv } (\text{Suc } n)) j)$  in asprod-0[of int N],
        assumption, simp add:less-ant-def, simp, simp,
        rule allI, rule impI, rule impI,
        rule Ostrowski-baseTr2, assumption+)
done

abbreviation
  CHOOSE :: [nat, nat]  $\Rightarrow$  nat  $(\langle \_ \rangle [80, 81] 80)$  where
    nCi == n choose i

```

lemma (in Ring) expansion-of-sum1: $x \in \text{carrier } R \implies$
 $(1_r \pm x)^{\sim R} n = \text{nsum } R (\lambda i. n C_i \times_R x^{\sim R} i) n$
apply (cut-tac ring-one, frule npeSum2[of 1_r x n], assumption+,
simp add:npOne, subgoal-tac $\forall (j::\text{nat}). (x^{\sim R} j) \in \text{carrier } R$)
apply (simp add:ring-l-one, rule allI, simp add:npClose)
done

lemma (in Ring) tail-of-expansion: $x \in \text{carrier } R \implies (1_r \pm x)^{\sim R} (\text{Suc } n) =$
 $(\text{nsum } R (\lambda i. ((\text{Suc } n) C_{(\text{Suc } i)} \times_R x^{\sim R} (\text{Suc } i))) n) \pm 1_r$
apply (cut-tac ring-is-ag)
apply (frule expansion-of-sum1[of x $\text{Suc } n$],
simp del:nsum-suc binomial-Suc-Suc npow-suc,
thin-tac $(1_r \pm x)^{\sim R} (\text{Suc } n) = \Sigma_e R (\lambda i. (\text{Suc } n) C_i \times_R x^{\sim R} i) (\text{Suc } n))$
apply (subst aGroup.nsumTail[of R n $\lambda i. (\text{Suc } n) C_i \times_R x^{\sim R} i$], assumption,
rule allI, rule impI, rule nsClose, rule npClose, assumption)
apply (cut-tac ring-one,
simp del:nsum-suc binomial-Suc-Suc npow-suc add:aGroup.ag-l-zero)
done

lemma (in Ring) tail-of-expansion1: $x \in \text{carrier } R \implies$
 $(1_r \pm x)^{\sim R} (\text{Suc } n) = x \cdot_r (\text{nsum } R (\lambda i. ((\text{Suc } n) C_{(\text{Suc } i)} \times_R x^{\sim R} i)) n) \pm 1_r$
apply (frule tail-of-expansion[of x n],
simp del:nsum-suc binomial-Suc-Suc npow-suc,
subgoal-tac $\forall i. \text{Suc } n C_{\text{Suc } i} \times_R x^{\sim R} i \in \text{carrier } R$,
cut-tac ring-one, cut-tac ring-is-ag)
prefer 2 **apply** (simp add: nsClose npClose)
apply (rule aGroup.ag-pOp-add-r[of R - 1_r], assumption+,
rule aGroup.nsum-mem, assumption+, rule allI, rule impI,
rule nsClose, rule npClose, assumption)
apply (rule ring-tOp-closed, assumption+,
rule aGroup.nsum-mem, assumption+, blast, simp add:ring-one)
apply (subst nsumMulEleL[of $\lambda i. \text{Suc } n C_{\text{Suc } i} \times_R x^{\sim R} i$ x], assumption+)
apply (rule aGroup.nsum-eq, assumption, rule allI, rule impI, rule nsClose,
rule npClose, assumption, rule allI, rule impI,
rule ring-tOp-closed, assumption+, rule nsClose, rule npClose,
assumption)
apply (rule allI, rule impI)
apply (subst nsMulDistrL, assumption, simp add:npClose,
frule-tac $n = j$ in npClose[of x], simp add:ring-tOp-commute[of x])
done

lemma (in Corps) nsum-in-VrTr: valuation K $v \implies$
 $(\forall j \leq n. f j \in \text{carrier } K) \wedge (\forall j \leq n.$
 $0 \leq (v (f j))) \longrightarrow (\text{nsum } K f n) \in \text{carrier } (Vr K v)$
apply (induct-tac n)
apply (rule impI, erule conjE, simp add:val-pos-mem-Vr)
apply (rule impI, erule conjE)
apply (frule Vr-ring[of v], frule Ring.ring-is-ag[of $Vr K v$],

$$\begin{aligned} & \text{frule-tac } x = f \text{ (Suc } n) \text{ and } y = \text{nsum } K \text{ f } n \text{ in} \\ & \quad \text{aGroup.ag-pOp-closed[of Vr K v],} \\ & \quad \text{subst val-pos-mem-Vr[THEN sym, of v], assumption+,} \\ & \quad \text{simp, simp, simp)} \\ \text{apply (simp, subst Vr-pOp-f-pOp[of v, THEN sym], assumption+,} \\ & \quad \text{subst val-pos-mem-Vr[THEN sym, of v], assumption+,} \\ & \quad \text{simp+)} \\ \text{apply (subst aGroup.ag-pOp-commute, assumption+, simp add:val-pos-mem-Vr,} \\ & \quad \text{assumption)} \\ \text{done} \end{aligned}$$

$$\begin{aligned} \text{lemma (in Corps) nsum-in-Vr:} & \llbracket \text{valuation } K \text{ v; } \forall j \leq n. f j \in \text{carrier } K; \\ & \forall j \leq n. 0 \leq (v (f j)) \rrbracket \implies (\text{nsum } K \text{ f } n) \in \text{carrier (Vr K v)} \\ \text{apply (simp add:nsum-in-VrTr)} \\ \text{done} \end{aligned}$$

$$\begin{aligned} \text{lemma (in Corps) nsum-mem-in-Vr:} & \llbracket \text{valuation } K \text{ v;} \\ & \forall j \leq n. (f j) \in \text{carrier } K; \forall j \leq n. 0 \leq (v (f j)) \rrbracket \implies \\ & (\text{nsum } K \text{ f } n) \in \text{carrier (Vr K v)} \\ \text{by (rule nsum-in-Vr)} \end{aligned}$$

$$\begin{aligned} \text{lemma (in Corps) val-nscale-ge-selfTr:} & \llbracket \text{valuation } K \text{ v; } x \in \text{carrier } K; 0 \leq v x \rrbracket \\ & \implies v x \leq v (n \times_K x) \\ \text{apply (cut-tac field-is-ring, induct-tac n, simp)} \\ \text{apply (simp add:value-of-zero,} \\ & \quad \text{simp,} \\ & \quad \text{frule-tac } y = n \times_K x \text{ in amin-le-plus[of v x], assumption+,} \\ & \quad \text{rule Ring.nsClose, assumption+)} \\ \text{apply (simp add:amin-def,} \\ & \quad \text{frule Ring.ring-is-ag[of K],} \\ & \quad \text{frule-tac } n = n \text{ in Ring.nsClose[of K x], assumption+,} \\ & \quad \text{simp add:aGroup.ag-pOp-commute)} \\ \text{done} \end{aligned}$$

$$\begin{aligned} \text{lemma (in Corps) ApproximationTr:} & \llbracket \text{valuation } K \text{ v; } x \in \text{carrier } K; 0 \leq (v x) \rrbracket \\ & \implies \\ & v x \leq (v (1_r \pm -_a ((1_r \pm x)^{\sim K} (\text{Suc } n)))) \\ \text{apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],} \\ & \quad \text{case-tac } x = \mathbf{0}_K, \\ & \quad \text{simp, frule Ring.ring-one[of K], simp add:aGroup.ag-r-zero,} \\ & \quad \text{simp add:Ring.npOne, simp add:Ring.ring-l-one, simp add:aGroup.ag-r-inv1,} \\ & \quad \text{subst Ring.tail-of-expansion1[of K x], assumption+,} \\ & \quad \text{frule Ring.ring-one[of K]}) \\ \text{apply (subgoal-tac (nsum } K \text{ (}\lambda i. \text{Suc } n C_{\text{Suc } i} \times_K x^{\sim K i}) n) \in \text{carrier (Vr K v),} \\ & \quad \text{frule Vr-mem-f-mem[of v (nsum } K \text{ (}\lambda i. \text{Suc } n C_{\text{Suc } i} \times_K x^{\sim K i}) n],} \\ & \quad \text{assumption+,} \\ & \quad \text{frule-tac } x = x \text{ and } y = \text{nsum } K \text{ (}\lambda i. \text{Suc } n C_{\text{Suc } i} \times_K x^{\sim K i}) n \text{ in} \\ & \quad \text{Ring.ring-tOp-closed[of K], assumption+,} \\ & \quad \text{subst aGroup.ag-pOp-commute[of K - 1_r], assumption+,} \end{aligned}$$

```

    subst aGroup.ag-p-inv[of K 1r], assumption+,
    subst aGroup.ag-pOp-assoc[THEN sym], assumption+,
    simp add:aGroup.ag-mOp-closed, rule aGroup.ag-mOp-closed, assumption+,
    simp del:binomial-Suc-Suc add:aGroup.ag-r-inv1, subst aGroup.ag-l-zero,
    assumption+,
    rule aGroup.ag-mOp-closed, assumption+, simp add:val-minus-eq)

apply (subst val-t2p[of v], assumption+) apply (
  simp add:val-pos-mem-Vr[THEN sym, of v
    nsum K (λi.(nCi + nCSuc i) ×K xK i) n],
  frule aadd-le-mono[of 0 v (nsum K (λi.(nCi + nCSuc i) ×K xK i) n)
    v x], simp add:aadd-0-l, simp add:aadd-commute[of v x])

apply (rule nsum-mem-in-Vr[of v n λi.Suc nCSuc i ×K xK i], assumption,
  rule allI, rule impI) apply (rule Ring.nsClose, assumption+) apply (simp
  add:Ring.npClose)

apply (rule allI, rule impI)
apply (cut-tac i = 0 and j = v (xK j) and k = v (Suc nCSuc j ×K xK j)
  in ale-trans)
apply (case-tac j = 0, simp add:value-of-one)
apply (simp add: val-exp-ring[THEN sym],
  frule val-nonzero-z[of v x], assumption+,
  erule exE,
  cut-tac m1 = 0 and n1 = j in of-nat-less-iff[THEN sym],
  frule-tac a = 0 < j and b = int 0 < int j in a-b-exchange,
  assumption, thin-tac 0 < j, thin-tac (0 < j) = (int 0 < int j))
apply (simp del: of-nat-0-less-iff)

apply (frule-tac w1 = int j and x1 = 0 and y1 = ant z in
  asprod-pos-mono[THEN sym],
  simp only:asprod-n-0)

apply(rule-tac x = xK j and n = Suc nCSuc j in
  val-nscal-ge-selfTr[of v], assumption+,
  simp add:Ring.npClose, simp add:val-exp-ring[THEN sym],
  frule val-nonzero-z[of v x], assumption+, erule exE, simp)
apply (case-tac j = 0, simp)
apply (subst asprod-amult, simp, simp add:a-z-z)
apply(
  simp only:ant-0[THEN sym], simp only:ale-zle,
  cut-tac m1 = 0 and n1 = j in of-nat-less-iff[THEN sym])
apply (
  frule-tac a = 0 < j and b = int 0 < int j in a-b-exchange,
  assumption+, thin-tac 0 < j, thin-tac (0 < j) = (int 0 < int j),
  frule-tac z = int 0 and z' = int j in zless-imp-zle,
  frule-tac i = int 0 and j = int j and k = z in int-mult-le,
  assumption+, simp add:mult.commute )
apply assumption
done

```

lemma (in Corps) *ApproximationTr0*: $aa \in \text{carrier } K \implies$
 $(1_r \pm -_a (aa \sim^K N)) \sim^K N \in \text{carrier } K$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
rule Ring.npClose, assumption+,
rule aGroup.ag-pOp-closed, assumption+, simp add:Ring.ring-one,
rule aGroup.ag-mOp-closed, assumption+, rule Ring.npClose, assumption+)
done

lemma (in Corps) *ApproximationTr1*: $aa \in \text{carrier } K \implies$
 $1_r \pm -_a ((1_r \pm -_a (aa \sim^K N)) \sim^K N) \in \text{carrier } K$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
frule ApproximationTr0[of aa N],
frule Ring.ring-one[of K], rule aGroup.ag-pOp-closed, assumption+,
rule aGroup.ag-mOp-closed, assumption+)
done

lemma (in Corps) *ApproximationTr2*: $\llbracket \text{valuation } K \ v; aa \in \text{carrier } K; aa \neq \mathbf{0};$
 $0 \leq (v \ aa) \rrbracket \implies (int \ N) *_a (v \ aa) \leq (v \ (1_r \pm -_a ((1_r \pm -_a (aa \sim^K N)) \sim^K N)))$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
case-tac $N = 0$,
frule val-nonzero-z[of v aa], assumption+, erule exE, simp)
apply (frule Ring.ring-one[of K], simp add:aGroup.ag-r-inv1,
simp add:value-of-zero)

apply (frule-tac $n = N$ in Ring.npClose[of K aa], assumption+,
frule ApproximationTr[of v -_a (aa \sim^K N) N - Suc 0],
rule aGroup.ag-mOp-closed, assumption+, simp add:val-minus-eq,
subst val-exp-ring[THEN sym, of v], assumption+,
simp add:asprod-pos-pos)
apply (simp add:val-minus-eq, simp add:val-exp-ring[THEN sym])
done

lemma (in Corps) *eSum-tr*:
 $(\forall j \leq n. (x \ j) \in \text{carrier } K) \wedge$
 $(\forall j \leq n. (b \ j) \in \text{carrier } K) \wedge l \leq n \wedge$
 $(\forall j \in (\{h. h \leq n\} - \{l\}). (g \ j = (x \ j) \cdot_r (1_r \pm -_a (b \ j)))) \wedge$
 $g \ l = (x \ l) \cdot_r (-_a (b \ l))$
 $\longrightarrow (nsum \ K \ (\lambda j \in \{h. h \leq n\}. (x \ j) \cdot_r (1_r \pm -_a (b \ j))) \ n) \pm (-_a (x \ l)) =$
 $nsum \ K \ g \ n$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (induct-tac n)
apply (simp, rule impI, (erule conjE)+,
simp, frule Ring.ring-one[of K], subst Ring.ring-distrib1,
assumption+,
simp add:aGroup.ag-mOp-closed, simp add:Ring.ring-r-one,
frule aGroup.ag-mOp-closed[of K b 0], assumption+,
frule Ring.ring-tOp-closed[of K x 0 -_a (b 0)], assumption+,
subst aGroup.ag-pOp-commute[of K x 0 -], assumption+,

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    subst aGroup.ag-pOp-assoc, assumption+,
    frule aGroup.ag-mOp-closed[of K],
    assumption+)
  apply (simp add:aGroup.ag-r-inv1, subst aGroup.ag-r-zero, assumption+, simp)
  apply (rule impI, (erule conjE)+)
  apply (subgoal-tac  $\forall j \leq (\text{Suc } n). ((x j) \cdot_r (1_r \pm -_a (b j))) \in \text{carrier } K$ )
  apply (case-tac  $l = \text{Suc } n$ , simp)
  apply (subgoal-tac  $\Sigma_e K g n \in \text{carrier } K$ ,
    subgoal-tac  $\{h. h \leq (\text{Suc } n)\} - \{\text{Suc } n\} = \{h. h \leq n\}$ , simp,
    subgoal-tac  $\forall j. j \leq n \longrightarrow j \leq (\text{Suc } n)$ ,
    frule-tac  $f = \lambda u. \text{if } u \leq \text{Suc } n \text{ then } (x u) \cdot_r (1_r \pm -_a (b u)) \text{ else undefined}$  and  $n = n$  in aGroup.nsum-eq[of K - - g])
  apply (rule allI, rule impI, simp,
    rule allI, simp, rule allI, rule impI, simp, simp)

  apply (cut-tac  $a = x (\text{Suc } n) \cdot_r (1_r \pm -_a (b (\text{Suc } n))) \pm -_a (x (\text{Suc } n))$  and
     $b = x (\text{Suc } n) \cdot_r (-_a (b (\text{Suc } n)))$  and
     $c = \Sigma_e K g n$  in aGroup.ag-pOp-add-l[of K], assumption)
  apply (rule aGroup.ag-pOp-closed, assumption+,
    rule Ring.ring-tOp-closed, assumption+, simp,
    rule aGroup.ag-pOp-closed, assumption+, simp add:Ring.ring-one,
    rule aGroup.ag-mOp-closed, assumption, simp,
    rule aGroup.ag-mOp-closed, assumption, simp,
    rule Ring.ring-tOp-closed, assumption+, simp,
    rule aGroup.ag-mOp-closed, assumption+, simp, assumption)

  apply (subst Ring.ring-distrib1, assumption+, simp, simp add:Ring.ring-one,
    simp add:aGroup.ag-mOp-closed,
    simp add:Ring.ring-r-one) apply (
    frule-tac  $x = x (\text{Suc } n)$  and  $y = x (\text{Suc } n) \cdot_r (-_a (b (\text{Suc } n)))$  in
    aGroup.ag-pOp-commute [of K], simp,
    simp add:Ring.ring-tOp-closed aGroup.ag-mOp-closed,
    simp) apply (
    subst aGroup.ag-pOp-assoc[of K], assumption+,
    rule Ring.ring-tOp-closed, assumption+, simp,
    (simp add:aGroup.ag-mOp-closed)+,
    subst aGroup.ag-r-inv1, assumption+, simp,
    subst aGroup.ag-r-zero, assumption+,
    simp add:Ring.ring-tOp-closed aGroup.ag-mOp-closed, simp,
    rotate-tac -1, drule sym, simp) apply (
    thin-tac  $\Sigma_e K g n \pm x (\text{Suc } n) \cdot_r (-_a (b (\text{Suc } n))) =$ 
     $\Sigma_e K g n \pm (x (\text{Suc } n) \cdot_r (1_r \pm -_a (b (\text{Suc } n))) \pm -_a (x (\text{Suc } n)))$ )
  apply (subst aGroup.ag-pOp-assoc[THEN sym], assumption+,
    rule Ring.ring-tOp-closed, assumption+, simp,
    rule aGroup.ag-pOp-closed, assumption+, simp add:Ring.ring-one,
    rule aGroup.ag-mOp-closed, assumption+, simp,
    rule aGroup.ag-mOp-closed, assumption+, simp, simp,
    simp, rule equalityI, rule subsetI, simp, rule subsetI, simp)
  apply (rule aGroup.nsum-mem, assumption+,

```



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      rule allI, rule impI, simp)
defer
  apply (rule allI, rule impI)
  apply (case-tac j = l, simp,
    rule Ring.ring-tOp-closed, assumption, simp,
    rule aGroup.ag-pOp-closed, assumption+, simp add:Ring.ring-one,
    rule aGroup.ag-mOp-closed, assumption, simp, simp,
    rule Ring.ring-tOp-closed, assumption, simp,
    rule aGroup.ag-pOp-closed, assumption+, simp add:Ring.ring-one,
    rule aGroup.ag-mOp-closed, assumption, simp, simp)

  apply (subst aGroup.ag-pOp-assoc, assumption+,
    rule aGroup.nsum-mem, assumption+,
    rule allI, simp, rule Ring.ring-tOp-closed, assumption+, simp,
    rule aGroup.ag-pOp-closed, assumption+, simp add:Ring.ring-one,
    rule aGroup.ag-mOp-closed, assumption, simp,
    rule aGroup.ag-mOp-closed, assumption, simp,
    subst aGroup.ag-pOp-commute[of K - -a (x l)], assumption+,
    rule Ring.ring-tOp-closed, assumption, simp,
    rule aGroup.ag-pOp-closed, assumption+, simp add:Ring.ring-one)
  apply (rule aGroup.ag-mOp-closed, assumption+, simp,
    rule aGroup.ag-mOp-closed, assumption+, simp,
    subst aGroup.ag-pOp-assoc[THEN sym], assumption+,
    rule aGroup.nsum-mem, assumption+,
    rule allI, rule impI, simp,
    rule aGroup.ag-mOp-closed, assumption, simp,
    rule Ring.ring-tOp-closed, assumption, simp,
    rule aGroup.ag-pOp-closed, assumption+, simp add:Ring.ring-one,
    rule aGroup.ag-mOp-closed, assumption, simp)
  apply (subgoal-tac  $\Sigma_e K (\lambda a. \text{if } a \leq (\text{Suc } n) \text{ then } x \ a \cdot_r (1_r \pm -_a (b \ a))$ 
    else undefined)  $n \pm -_a (x \ l) =$ 
 $\Sigma_e K (\lambda a. \text{if } a \leq n \text{ then } x \ a \cdot_r (1_r \pm -_a (b \ a))$  else undefined)  $n \pm$ 
 $-_a (x \ l)$ , simp,
    rule aGroup.ag-pOp-add-r[of K - -a (x l)], assumption+,
    rule aGroup.nsum-mem, assumption+,
    rule allI, rule impI, simp,
    rule aGroup.nsum-mem, assumption+,
    rule allI, rule impI, simp,
    rule aGroup.ag-mOp-closed, assumption, simp,
    rule aGroup.nsum-eq, assumption+,
    rule allI, rule impI, simp, rule allI, rule impI)
  apply simp
  apply (rule allI, rule impI, simp)
done

```

lemma (in Corps) *eSum-minus-x*: $\llbracket \forall j \leq n. (x \ j) \in \text{carrier } K;$
 $\forall j \leq n. (b \ j) \in \text{carrier } K; l \leq n;$
 $\forall j \in (\{h. h \leq n\} - \{l\}). (g \ j = (x \ j) \cdot_r (1_r \pm -_a (b \ j)))$;
 $g \ l = (x \ l) \cdot_r (-_a (b \ l)) \rrbracket \implies$

$$(nsum\ K\ (\lambda j \in \{h. h \leq n\}. (x\ j) \cdot_r (1_r \pm -_a (b\ j)))\ n) \pm (-_a (x\ l)) =$$

$$nsum\ K\ g\ n$$
by (*cut-tac eSum-tr*[*of* $n\ x\ b\ l\ g$], *simp*)

lemma (**in** *Ring*) *one-m-x-times*: $x \in carrier\ R \implies$
 $(1_r \pm -_a x) \cdot_r (nsum\ R\ (\lambda j. x^{\sim R} j)\ n) = 1_r \pm -_a (x^{\sim R} (Suc\ n))$
apply (*cut-tac ring-one*, *cut-tac ring-is-ag*,
frule aGroup.ag-mOp-closed[*of* $R\ x$], *assumption+*,
frule aGroup.ag-pOp-closed[*of* $1_r -_a x$], *assumption+*)

apply (*induct-tac n*)
apply (*simp add:ring-r-one ring-l-one*)
apply (*simp del:npow-suc*,
frule-tac n = Suc n in npClose[*of* x],
subst ring-distrib1, *assumption+*)
apply (*rule aGroup.nsum-mem*, *assumption*, *rule allI*, *rule impI*,
simp add:npClose, *rule npClose*, *assumption+*,
simp del:npow-suc,
thin-tac ($1_r \pm -_a x) \cdot_r \Sigma_e R\ (npow\ R\ x)\ n = 1_r \pm -_a (x^{\sim R} (Suc\ n))$)
apply (*subst ring-distrib2*, *assumption+*,
simp del:npow-suc add:ring-l-one,
subst aGroup.pOp-assocTr43[*of* R], *assumption+*,
rule-tac $x = x^{\sim R} (Suc\ n)$ **in** *aGroup.ag-mOp-closed*[*of* R], *assumption+*,
rule ring-tOp-closed, *rule aGroup.ag-mOp-closed*, *assumption+*)
apply (*subst aGroup.ag-l-inv1*, *assumption+*, *simp del:npow-suc*
add:aGroup.ag-r-zero,
frule-tac $x = -_a x$ **and** $y = x^{\sim R} (Suc\ n)$ **in** *ring-tOp-closed*,
assumption+)
apply (*rule aGroup.ag-pOp-add-l*[*of* $R - 1_r$], *assumption+*,
rule aGroup.ag-mOp-closed, *assumption+*,
rule npClose, *assumption+*,
subst ring-inv1-1[*THEN sym*, *of* x], *assumption*,
rule npClose, *assumption*,
simp,
subst ring-tOp-commute[*of* x], *assumption+*, *simp*)

done

lemma (**in** *Corps*) *x-pow-fSum-in-Vr*: $\llbracket valuation\ K\ v; x \in carrier\ (Vr\ K\ v) \rrbracket \implies$
 $(nsum\ K\ (npow\ K\ x)\ n) \in carrier\ (Vr\ K\ v)$
apply (*frule Vr-ring*[*of* v])
apply (*induct-tac n*)
apply *simp*
apply (*frule Ring.ring-one*[*of* $Vr\ K\ v$])
apply (*simp add:Vr-1-f-1*)
apply (*simp del:npow-suc*)
apply (*frule Ring.ring-is-ag*[*of* $Vr\ K\ v$])

apply (*subst Vr-pOp-f-pOp*[*THEN sym*, *of* v], *assumption+*)
apply (*subst Vr-exp-f-exp*[*THEN sym*, *of* v], *assumption+*)

```

apply (rule Ring.npClose[of Vr K v], assumption+)
apply (rule aGroup.ag-pOp-closed[of Vr K v], assumption+)
apply (subst Vr-exp-f-exp[THEN sym, of v], assumption+)
apply (rule Ring.npClose[of Vr K v], assumption+)
done

lemma (in Corps) val-1mx-pos: [valuation K v; x ∈ carrier K;
  0 < (v (1r ± -a x))] ⇒ v x = 0
apply (cut-tac field-is-ring, frule Ring.ring-one[of K],
  frule Ring.ring-is-ag[of K])
apply (frule aGroup.ag-mOp-closed[of K x], assumption+)
apply (frule aGroup.ag-pOp-closed[of K 1r -a x], assumption+)
apply (frule aGroup.ag-mOp-closed[of K 1r ± -a x], assumption+)
apply (cut-tac x = x and y = 1r ± -a (1r ± -a x) and f = v in
  eq-elems-eq-val)
apply (subst aGroup.ag-p-inv, assumption+,
  subst aGroup.ag-pOp-assoc[THEN sym], assumption+,
  rule aGroup.ag-mOp-closed, assumption+,
  subst aGroup.ag-inv-inv, assumption+,
  subst aGroup.ag-r-inv1, assumption+,
  subst aGroup.ag-l-zero, assumption+,
  (simp add: aGroup.ag-inv-inv)+,
  frule value-less-eq[of v 1r -a (1r ± -a x)],
  assumption+)
apply (simp add: val-minus-eq value-of-one,
  simp add: value-of-one)
done

lemma (in Corps) val-1mx-pow: [valuation K v; x ∈ carrier K;
  0 < (v (1r ± -a x))] ⇒ 0 < (v (1r ± -a xK (Suc n)))
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (subst Ring.one-m-x-times[THEN sym, of K x n], assumption+)
apply (frule Ring.ring-one[of K],
  frule x-pow-fSum-in-Vr[of v x n],
  subst val-pos-mem-Vr[THEN sym], assumption+,
  frule val-1mx-pos[of v x], assumption+,
  simp)

apply (subst val-t2p, assumption+,
  rule aGroup.ag-pOp-closed, assumption+,
  simp add: aGroup.ag-mOp-closed, simp add: Vr-mem-f-mem,
  frule val-pos-mem-Vr[THEN sym, of v nsum K (npow K x) n],
  simp add: Vr-mem-f-mem, simp)
apply (frule aadd-le-mono[of 0 v (nsum K (npow K x) n) v (1r ± -a x)],
  simp add: aadd-0-l, simp add: aadd-commute)
done

lemma (in Corps) ApproximationTr3: [vals-nonequiv K (Suc n) vv;
  ∀ l ≤ (Suc n). x l ∈ carrier K; j ≤ (Suc n)] ⇒

```

$\exists L. (\forall N. L < N \longrightarrow (an\ m) \leq (vv\ j\ ((\Sigma_e\ K\ (\lambda k \in \{h. h \leq (Suc\ n)\}. (x\ k) \cdot_r (1_r \pm -_a ((1_r \pm -_a (((\Omega_K\ vv\ (Suc\ n))\ k) \sim^K N)) \sim^K N))) (Suc\ n)) \pm -_a (x\ j))))$
apply (*cut-tac field-is-ring, frule Ring.ring-is-ag*[of K])
apply (*frule-tac vals-nonequiv-valuation*[of $Suc\ n\ vv\ j$], *assumption*+)
apply (*subgoal-tac* $\forall N. \Sigma_e\ K\ (\lambda j \in \{h. h \leq (Suc\ n)\}. (x\ j) \cdot_r (1_r \pm -_a (1_r \pm -_a ((\Omega_K\ vv\ (Suc\ n))\ j) \sim^K N) \sim^K N)) (Suc\ n) \pm -_a (x\ j) = \Sigma_e\ K\ (\lambda l \in \{h. h \leq (Suc\ n)\}. (if\ l \neq j\ then\ (x\ l) \cdot_r (1_r \pm -_a (1_r \pm -_a ((\Omega_K\ vv\ (Suc\ n))\ l) \sim^K N) \sim^K N) \ else\ (x\ j) \cdot_r (1_r \pm -_a (1_r \pm -_a ((\Omega_K\ vv\ (Suc\ n))\ l) \sim^K N) \sim^K N \pm -_a\ 1_r))) (Suc\ n))$)
apply (*simp del:nsum-suc*)
apply (*thin-tac* $\forall N. \Sigma_e\ K\ (\lambda j \in \{h. h \leq (Suc\ n)\}. (x\ j) \cdot_r (1_r \pm -_a (1_r \pm -_a ((\Omega_K\ vv\ (Suc\ n))\ j) \sim^K N) \sim^K N)) (Suc\ n) \pm -_a (x\ j) = \Sigma_e\ K\ (\lambda l \in \{h. h \leq (Suc\ n)\}. (if\ l \neq j\ then\ (x\ l) \cdot_r (1_r \pm -_a (1_r \pm -_a ((\Omega_K\ vv\ (Suc\ n))\ l) \sim^K N) \sim^K N) \ else\ (x\ j) \cdot_r (1_r \pm -_a (1_r \pm -_a ((\Omega_K\ vv\ (Suc\ n))\ l) \sim^K N) \sim^K N \pm -_a\ 1_r))) (Suc\ n))$)
prefer 2 apply (*rule allI*)
apply (*rule eSum-minus-x, assumption*+)
apply (*rule allI, rule impI*) **apply** (*rule ApproximationTr0*)
apply (*simp add:Ostrowski-base-mem*) **apply** *assumption*
apply (*rule ballI, simp*)
apply *simp*
apply (*frule Ring.ring-one*[of K])
apply (*cut-tac* $aa = (\Omega_K\ vv\ (Suc\ n))\ j\ and\ N = N\ in$
ApproximationTr0)
apply (*simp add:Ostrowski-base-mem*)
apply (*subst aGroup.ag-pOp-assoc, assumption*+)
apply (*rule aGroup.ag-mOp-closed, assumption*+) +
apply (*subst aGroup.ag-pOp-commute*[of $K - -_a\ 1_r$], *assumption*+) +
apply (*rule aGroup.ag-mOp-closed, assumption*+) +
apply (*subst aGroup.ag-pOp-assoc*[*THEN sym*], *assumption*+) +
apply (*rule aGroup.ag-mOp-closed, assumption*+) +
apply (*simp add:aGroup.ag-r-inv1*)
apply (*subst aGroup.ag-l-zero, assumption*+) **apply** (*simp add:aGroup.ag-mOp-closed*)
apply *simp*
apply (*subgoal-tac* $\exists L. \forall N. L < N \longrightarrow$
 $(\forall ja \leq (Suc\ n). (an\ m) \leq ((vv\ j \circ (\lambda l \in \{h. h \leq (Suc\ n)\}. (if\ l \neq j\ then\ (x\ l) \cdot_r (1_r \pm -_a (1_r \pm -_a ((\Omega_K\ vv\ (Suc\ n))\ l) \sim^K N) \sim^K N) \ else\ (x\ j) \cdot_r (1_r \pm -_a (1_r \pm -_a ((\Omega_K\ vv\ (Suc\ n))\ l) \sim^K N) \sim^K N \pm -_a\ 1_r))) ja)))$)
apply (*erule exE*)
apply (*rename-tac M*)
apply (*subgoal-tac* $\forall N. M < (N::nat) \longrightarrow$

$(an\ m) \leq (vv\ j\ (\Sigma_e\ K\ (\lambda l \in \{h.\ h \leq (Suc\ n)\}).\ (if\ l \neq j\ then$
 $(x\ l) \cdot_r\ (1_r \pm -_a\ (1_r \pm -_a\ ((\Omega_K\ vv\ (Suc\ n))\ l) \neg^K N) \neg^K N)$
 $else\ (x\ j) \cdot_r\ (1_r \pm -_a\ (1_r \pm -_a\ ((\Omega_K\ vv\ (Suc\ n))\ l) \neg^K N) \neg^K N$
 $\pm -_a\ 1_r)))\ (Suc\ n))))$
apply *blast*
apply (rule *allI*, rule *impI*)
apply (drule-tac *a = N in forall-spec*, *assumption*)
apply (rule *value-ge-add*[of *vv j Suc n - an m*], *assumption+*)

apply (rule *allI*, rule *impI*)
apply (frule *Ring.ring-one*[of *K*])
apply (case-tac *ja = j*, *simp*)
apply (rule *Ring.ring-tOp-closed*, *assumption+*, *simp*)
apply (rule *aGroup.ag-pOp-closed*, *assumption+*)
apply (rule *aGroup.ag-mOp-closed*, *assumption+*)
apply (rule *Ring.npClose*, *assumption*)
apply (rule *aGroup.ag-pOp-closed*, *assumption+*)
apply (rule *aGroup.ag-mOp-closed*, *assumption*)
apply (rule *Ring.npClose*, *assumption*)
apply (simp *add:Ostrowski-base-mem*)
apply (rule *aGroup.ag-mOp-closed*, *assumption+*)

apply *simp*
apply (rule *Ring.ring-tOp-closed*, *assumption+*, *simp*)
apply (rule *aGroup.ag-pOp-closed*, *assumption+*)
apply (rule *aGroup.ag-mOp-closed*, *assumption+*)
apply (rule *Ring.npClose*, *assumption*)
apply (rule *aGroup.ag-pOp-closed*, *assumption+*)
apply (rule *aGroup.ag-mOp-closed*, *assumption*)
apply (rule *Ring.npClose*, *assumption*)
apply (simp *add:Ostrowski-base-mem*)

apply *assumption*

apply (subgoal-tac $\forall N. \forall ja \leq (Suc\ n). (1_r \pm -_a\ (1_r \pm -_a$
 $((\Omega_K\ vv\ (Suc\ n))\ ja) \neg^K N) \neg^K N) \in carrier\ K)$
apply (subgoal-tac $\forall N. (1_r \pm -_a\ (1_r \pm -_a\ ((\Omega_K\ vv\ (Suc\ n))\ j) \neg^K N) \neg^K N$
 $\pm -_a\ 1_r) \in carrier\ K)$
apply (simp *add:val-t2p*)
apply (cut-tac *multi-inequalityTr0*[of *Suc n (vv j) o x m*])
apply (subgoal-tac $\forall ja \leq (Suc\ n). (vv\ j \circ x)\ ja \neq -\infty$, *simp*)
apply (erule *exE*)
apply (subgoal-tac $\forall N. L < N \longrightarrow (\forall ja \leq (Suc\ n). (ja \neq j \longrightarrow$
 $an\ m \leq vv\ j\ (x\ ja) + (vv\ j\ (1_r \pm -_a\ (1_r \pm -_a\ ((\Omega_K\ vv\ (Suc\ n))\ ja) \neg^K N) \neg^K N)))$
 $\wedge (ja = j \longrightarrow (an\ m) \leq vv\ j\ (x\ j) + (vv\ j\ (1_r \pm -_a\ (1_r \pm$
 $-_a\ ((\Omega_K\ vv\ (Suc\ n))\ j) \neg^K N) \neg^K N \pm -_a\ (1_r))))))$
apply *blast*

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apply (rule allI, rule impI)+

apply (case-tac ja = j, simp)
apply (thin-tac  $\forall N. 1_r \pm -_a (1_r \pm -_a (\Omega_K \text{ vv } (\text{Suc } n)) j \sim^K N) \sim^K N \pm -_a 1_r$ 
 $\in$ 
  carrier K)
apply (thin-tac  $\forall l \leq \text{Suc } n. x \ l \in \text{carrier } K$ )
apply (drule-tac  $x = N$  in spec)
apply (drule-tac  $a = j$  in forall-spec, assumption,
  thin-tac  $\forall ja \leq \text{Suc } n. 1_r \pm -_a (1_r \pm -_a (\Omega_K \text{ vv } (\text{Suc } n)) ja \sim^K N) \sim^K N$ 
 $\in$  carrier K)
apply (cut-tac  $N = N$  in ApproximationTr0 [of  $(\Omega_K \text{ vv } (\text{Suc } n)) j$ ])
apply (simp add:Ostrowski-base-mem)
apply (frule Ring.ring-one[of K], frule aGroup.ag-mOp-closed[of K 1_r],
  assumption) apply (
  frule-tac  $x = (1_r \pm -_a ((\Omega_K \text{ vv } (\text{Suc } n)) j) \sim^K N) \sim^K N$  in
  aGroup.ag-mOp-closed[of K], assumption+)
apply (simp only:aGroup.ag-pOp-assoc)
apply (simp only:aGroup.ag-pOp-commute[of K -  $-_a 1_r$ ])
apply (simp only:aGroup.ag-pOp-assoc[THEN sym])
apply (simp add:aGroup.ag-r-inv1)
apply (simp add:aGroup.ag-l-zero) apply (simp only:val-minus-eq)
apply (thin-tac  $(1_r \pm -_a (\Omega_K \text{ vv } (\text{Suc } n)) j \sim^K N) \sim^K N \in \text{carrier } K$ ,
  thin-tac  $-_a (1_r \pm -_a (\Omega_K \text{ vv } (\text{Suc } n)) j \sim^K N) \sim^K N \in \text{carrier } K$ )
apply (subst val-exp-ring[THEN sym, of vv j], assumption+)
apply (rule aGroup.ag-pOp-closed[of K], assumption+)
apply (rule aGroup.ag-mOp-closed[of K], assumption)
apply (rule Ring.npClose, assumption+) apply (simp add:Ostrowski-base-mem)
apply (rule Ostrowski-base-npow-not-one) apply simp apply assumption+
apply (drule-tac  $a = N$  in forall-spec, assumption)
apply (drule-tac  $a = j$  in forall-spec, assumption)
apply (frule Ostrowski-baseTr1[of n vv j], assumption+)
apply (frule-tac  $n = N - \text{Suc } 0$  in val-1mx-pow[of vv j  $(\Omega_K \text{ vv } (\text{Suc } n)) j$ ])
apply (simp add:Ostrowski-base-mem) apply assumption
apply (thin-tac  $\text{vv } j (x j) \neq -\infty$ ) apply (simp only:Suc-pred)
apply (thin-tac  $0 < \text{vv } j (1_r \pm -_a ((\Omega_K \text{ vv } (\text{Suc } n)) j)))$ 
apply (cut-tac  $b = \text{vv } j (1_r \pm -_a ((\Omega_K \text{ vv } (\text{Suc } n)) j) \sim^K N)$  and  $N = N$  in
  asprod-ge) apply assumption apply simp
apply (cut-tac  $x = \text{an } N$  and  $y = \text{int } N *_a \text{vv } j (1_r \pm -_a ((\Omega_K \text{ vv } (\text{Suc } n))$ 
 $j) \sim^K N)$  in aadd-le-mono[of - - vv j (x j)], assumption)
apply (simp add:aadd-commute)

apply simp
apply (frule-tac  $aa = (\Omega_K \text{ vv } (\text{Suc } n)) ja$  and  $N = N$  in
  ApproximationTr2[of vv j])
apply (simp add:Ostrowski-base-mem)

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    apply (rule Ostrowski-base-nonzero, assumption+)
  apply (frule-tac l = ja in Ostrowski-baseTr0[of n vv], assumption+,
    erule conjE)
  apply (rotate-tac -1, frule-tac a = j in forall-spec) apply assumption
  apply (frule-tac x = j in bspec, simp)
  apply (rule aless-imp-le) apply blast
  apply (rotate-tac -5,
    drule-tac a = N in forall-spec, assumption)
  apply (rotate-tac -2,
    drule-tac a = ja in forall-spec, assumption) apply (
    drule-tac a = ja in forall-spec, assumption)
  apply (frule-tac l = ja in Ostrowski-baseTr0[of n vv], assumption+)
  apply (erule conjE, rotate-tac -1,
    frule-tac a = j in forall-spec, assumption+)
  apply (thin-tac vv j (x ja)  $\neq -\infty$ )
  apply (cut-tac b = vv j (( $\Omega_K$  vv (Suc n)) ja) and N = N in asprod-ge)
  apply simp apply simp
  apply (frule-tac x = an N and y = int N *a vv j (( $\Omega_K$  vv (Suc n)) ja) and
    z = vv j (x ja) in aadd-le-mono)
  apply (frule-tac x = int N *a vv j (( $\Omega_K$  vv (Suc n)) ja) and y = (vv j)
    ( $1_r \pm -_a (1_r \pm -_a ((\Omega_K$  vv (Suc n)) ja)  $\neg^K N) \neg^K N$ ) and z = vv j (x ja)
    in aadd-le-mono)
  apply (frule-tac i = an N + vv j (x ja) and
    j = int N *a vv j (( $\Omega_K$  vv (Suc n)) ja) + vv j (x ja) and
    k = vv j ( $1_r \pm -_a (1_r \pm -_a ((\Omega_K$  vv (Suc n)) ja)  $\neg^K N) \neg^K N$ ) +
    vv j (x ja) in ale-trans, assumption+)
  apply (subst aadd-commute)
  apply (frule-tac x = an m and y = vv j (x ja) + an N in aless-imp-le)
  apply (rule-tac j = vv j (x ja) + an N in ale-trans[of an m],
    assumption)
  apply (simp add:aadd-commute)
  apply (rule allI, rule impI, subst comp-def)
  apply (frule-tac a = ja in forall-spec, assumption)
  apply (frule-tac x = x ja in value-in-aug-inf[of vv j], assumption+)
  apply (simp add:aug-inf-def)

  apply (rule allI)
  apply (rule aGroup.ag-pOp-closed, assumption+) apply blast
  apply (rule aGroup.ag-mOp-closed, assumption, rule Ring.ring-one, assumption)

  apply ((rule allI)+, rule impI)
  apply (rule-tac aa = ( $\Omega_K$  vv (Suc n)) ja in ApproximationTr1,
    simp add:Ostrowski-base-mem)
done

definition
  app-lb :: [- , nat, nat  $\Rightarrow$  'b  $\Rightarrow$  ant, nat  $\Rightarrow$  'b, nat]  $\Rightarrow$ 
    (nat  $\Rightarrow$  nat) ( $\langle (5\Psi$  - - - )  $\rangle$  [98,98,98,98,99]98) where

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$\Psi_K n \text{ vv } x m = (\lambda j \in \{h. h \leq n\}. (\text{SOME } L. (\forall N. L < N \longrightarrow$
 $(an \ m) \leq (vv \ j \ (\Sigma_e \ K \ (\lambda j \in \{h. h \leq n\}. (x \ j) \cdot_r \ K \ (1_r \ K \pm_K -_a \ K$
 $(1_r \ K \pm_K -_a \ K \ ((\Omega_K \text{ vv } n) \ j) \sim^K N) \sim^K N)) \ n \pm_K -_a \ K \ (x \ j))))))$

lemma (in Corps) app-LB: $\llbracket \text{vals-nonequiv } K \ (Suc \ n) \ \text{vv};$
 $\forall l \leq (Suc \ n). \ x \ l \in \text{carrier } K; \ j \leq (Suc \ n) \rrbracket \implies$
 $\forall N. (\Psi_K (Suc \ n) \ \text{vv } x \ m) \ j < N \longrightarrow (an \ m) \leq$
 $(vv \ j \ (\Sigma_e \ K \ (\lambda j \in \{h. h \leq (Suc \ n)\}. (x \ j) \cdot_r \ (1_r \pm -_a \ (1_r \pm$
 $-_a \ ((\Omega_K \text{ vv } (Suc \ n)) \ j) \sim^K N) \sim^K N)) \ (Suc \ n) \pm -_a \ (x \ j)))$
apply (frule ApproximationTr3[of n vv x j m],
 assumption+)
apply (simp del:nsum-suc add:app-lb-def) **apply** (rule allI)
apply (rule someI2-ex) **apply** assumption+
apply (rule impI) **apply** blast
done

lemma (in Corps) ApplicationTr4: $\llbracket \text{vals-nonequiv } K \ (Suc \ n) \ \text{vv};$
 $\forall j \in \{h. h \leq (Suc \ n)\}. \ x \ j \in \text{carrier } K \rrbracket \implies$
 $\exists l. \forall N. \ l < N \longrightarrow (\forall j \leq (Suc \ n). \ (an \ m) \leq$
 $(vv \ j \ (\Sigma_e \ K \ (\lambda j \in \{h. h \leq (Suc \ n)\}. (x \ j) \cdot_r \ (1_r \pm -_a \ (1_r \pm$
 $-_a \ ((\Omega_K \text{ vv } (Suc \ n)) \ j) \sim^K N) \sim^K N)) \ (Suc \ n) \pm -_a \ (x \ j))))$
apply (subgoal-tac $\forall N. \ (m\text{-max } (Suc \ n) \ (\Psi_K (Suc \ n) \ \text{vv } x \ m)) < N \longrightarrow$
 $(\forall j \leq (Suc \ n). \ (an \ m) \leq$
 $(vv \ j \ (\Sigma_e \ K \ (\lambda j \in \{h. h \leq (Suc \ n)\}. (x \ j) \cdot_r \ (1_r \pm -_a \ (1_r \pm$
 $-_a \ ((\Omega_K \text{ vv } (Suc \ n)) \ j) \sim^K N) \sim^K N)) \ (Suc \ n) \pm -_a \ (x \ j))))$)
apply blast
apply (rule allI, rule impI)+
apply (frule-tac $j = j$ in app-LB[of n vv x - m],
 $\text{simp}, \text{assumption},$
 $\text{subgoal-tac } (\Psi_K (Suc \ n) \ \text{vv } x \ m) \ j < N, \text{blast})$)
apply (frule-tac $l = j$ and $n = Suc \ n$ and $f = \Psi_K (Suc \ n) \ \text{vv } x \ m$ in m-max-gt,
 $\text{frule-tac } x = (\Psi_K (Suc \ n) \ \text{vv } x \ m) \ j$ and
 $y = m\text{-max } (Suc \ n) \ (\Psi_K (Suc \ n) \ \text{vv } x \ m)$ and $z = N$ in le-less-trans,
 assumption+)
done

theorem (in Corps) Approximation-thm: $\llbracket \text{vals-nonequiv } K \ (Suc \ n) \ \text{vv};$
 $\forall j \leq (Suc \ n). \ (x \ j) \in \text{carrier } K \rrbracket \implies$
 $\exists y \in \text{carrier } K. \forall j \leq (Suc \ n). \ (an \ m) \leq (vv \ j \ (y \pm -_a \ (x \ j)))$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (subgoal-tac $\exists l. (\forall N. \ l < N \longrightarrow (\forall j \leq (Suc \ n). \ (an \ m) \leq ((vv \ j) ((nsum \ K$
 $(\lambda j \in \{h. h \leq (Suc \ n)\}. (x \ j) \cdot_r \ (1_r \pm -_a \ (1_r \pm -_a \ ((\Omega_K \text{ vv } (Suc \ n)) \ j) \sim^K N) \sim^K N))$
 $(Suc \ n)) \pm -_a \ (x \ j))))))$)
apply (erule exE)
apply (rename-tac M)
apply (subgoal-tac $\forall j \leq (Suc \ n). \ (an \ m) \leq$

$(vv\ j\ (\ (\Sigma_e\ K\ (\lambda j \in \{h. h \leq (Suc\ n)\}). (x\ j) \cdot_r (1_r \pm -_a (1_r \pm$
 $-_a ((\Omega_K\ vv\ (Suc\ n))\ j) \sim^K (Suc\ M)) \sim^K (Suc\ M))) (Suc\ n)) \pm -_a (x\ j))))$
apply (subgoal-tac $\Sigma_e\ K\ (\lambda j \in \{h. h \leq (Suc\ n)\}). (x\ j) \cdot_r (1_r \pm$
 $-_a (1_r \pm -_a ((\Omega_K\ vv\ (Suc\ n))\ j) \sim^K (Suc\ M)) \sim^K (Suc\ M))) (Suc\ n) \in carrier\ K)$
apply blast
apply (rule aGroup.nsum-mem[of $K\ Suc\ n$], assumption+)
apply (rule allI, rule impI, simp del:nsum-suc npow-suc)
apply (rule Ring.ring-tOp-closed, assumption+, simp,
rule ApproximationTr1, simp add:Ostrowski-base-mem)

apply (subgoal-tac $M < Suc\ M$) **apply** blast
apply simp
apply (rule ApplicationTr4[of $n\ vv\ x$], assumption+)
apply simp
done

definition

distinct-pds :: $[-, nat, nat \Rightarrow ('b \Rightarrow ant)\ set] \Rightarrow bool$ **where**
distinct-pds $K\ n\ P \longleftrightarrow (\forall j \leq n. P\ j \in Pds\ K) \wedge$
 $(\forall l \leq n. \forall m \leq n. l \neq m \longrightarrow P\ l \neq P\ m)$

lemma (in Corps) distinct-pds-restriction: $\llbracket distinct-pds\ K\ (Suc\ n)\ P \rrbracket \Longrightarrow$
distinct-pds $K\ n\ P$
apply (simp add:distinct-pds-def)
done

lemma (in Corps) ring-n-distinct-prime-divisors: distinct-pds $K\ n\ P \Longrightarrow$
Ring (Sr $K\ \{x. x \in carrier\ K \wedge (\forall j \leq n. 0 \leq ((\nu_K\ (P\ j))\ x))\}$)
apply (simp add:distinct-pds-def) **apply** (erule conjE)
apply (cut-tac field-is-ring)
apply (rule Ring.Sr-ring, assumption+)
apply (subst sr-def)
apply (rule conjI)
apply (rule subsetI) **apply** simp
apply (rule conjI)
apply (simp add:Ring.ring-one)
apply (rule allI, rule impI)
apply (cut-tac $P = P\ j$ in representative-of-pd-valuation, simp,
simp add:value-of-one)
apply (rule ballI)+
apply simp
apply (frule Ring.ring-is-ag[of K]) **apply** (erule conjE)+
apply (frule-tac $x = y$ in aGroup.ag-mOp-closed[of K], assumption+)
apply (frule-tac $x = x$ and $y = -_a\ y$ in aGroup.ag-pOp-closed[of K],
assumption+)
apply simp
apply (rule conjI)

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apply (rule allI, rule impI)
apply (rotate-tac -4, frule-tac a = j in forall-spec, assumption,
        rotate-tac -3,
        drule-tac a = j in forall-spec, assumption)
apply (cut-tac P = P j in representative-of-pd-valuation, simp)
apply (frule-tac v =  $\nu_K$  (P j) and x = x and y =  $-\text{a}$  y in amin-le-plus,
        assumption+)
apply (simp add:val-minus-eq)
apply (frule-tac x = ( $\nu_K$  (P j)) x and y = ( $\nu_K$  (P j)) y in amin-ge1[of 0])
        apply simp
apply (rule-tac j = amin (( $\nu_K$  (P j)) x) (( $\nu_K$  (P j)) y) and k = ( $\nu_K$  (P j)) (x  $\pm$ 
 $-\text{a}$  y) in ale-trans[of 0], assumption+)
apply (simp add:Ring.ring-tOp-closed)

apply (rule allI, rule impI,
        cut-tac P = P j in representative-of-pd-valuation, simp,
        subst val-t2p [where v= $\nu_K$  P j], assumption+,
        rule aadd-two-pos, simp+)
done

lemma (in Corps) distinct-pds-valuation:  $\llbracket j \leq (\text{Suc } n);$ 
        distinct-pds K (Suc n) P  $\rrbracket \implies$  valuation K ( $\nu_K$  (P j))
apply (rule-tac P = P j in representative-of-pd-valuation)
apply (simp add:distinct-pds-def)
done

lemma (in Corps) distinct-pds-valuation1:  $\llbracket 0 < n; j \leq n; \text{distinct-pds K } n \text{ P} \rrbracket$ 
 $\implies$  valuation K ( $\nu_K$  (P j))
apply (rule distinct-pds-valuation[of j n - Suc 0 P])
apply simp+
done

lemma (in Corps) distinct-pds-valuation2:  $\llbracket j \leq n; \text{distinct-pds K } n \text{ P} \rrbracket \implies$ 
        valuation K ( $\nu_K$  (P j))
apply (case-tac n = 0,
        simp add:distinct-pds-def,
        subgoal-tac 0  $\in \{0::\text{nat}\}$ ,
        simp add:representative-of-pd-valuation[of P 0],
        simp)

apply (simp add:distinct-pds-valuation1[of n])
done

definition
ring-n-pd ::  $\llbracket ('b, 'm) \text{ Ring-scheme}, \text{nat} \Rightarrow ('b \Rightarrow \text{ant}) \text{ set},$ 
        nat  $\rrbracket \Rightarrow ('b, 'm) \text{ Ring-scheme}$ 
        ( $\langle \langle \exists O. \dots \rangle, [98, 98, 99] 98 \rangle$ ) where
 $O_K P n = \text{Sr K } \{x. x \in \text{carrier K} \wedge$ 

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$$(\forall j \leq n. 0 \leq ((\nu_K (P j)) x))\}$$

lemma (in Corps) ring-n-pd:distinct-pds $K \ n \ P \implies \text{Ring } (O_K \ P \ n)$
by (simp add:ring-n-pd-def, simp add:ring-n-distinct-prime-divisors)

lemma (in Corps) ring-n-pd-Suc:distinct-pds $K \ (\text{Suc } n) \ P \implies$
 $\text{carrier } (O_K \ P \ (\text{Suc } n)) \subseteq \text{carrier } (O_K \ P \ n)$
apply (rule subsetI)
apply (simp add:ring-n-pd-def Sr-def)
done

lemma (in Corps) ring-n-pd-pOp-K-pOp: $\llbracket \text{distinct-pds } K \ n \ P; x \in \text{carrier } (O_K \ P \ n);$
 $y \in \text{carrier } (O_K \ P \ n) \rrbracket \implies x \pm_{(O_K \ P \ n)} y = x \pm y$
apply (simp add:ring-n-pd-def Sr-def)
done

lemma (in Corps) ring-n-pd-tOp-K-tOp: $\llbracket \text{distinct-pds } K \ n \ P; x \in \text{carrier } (O_K \ P \ n);$
 $y \in \text{carrier } (O_K \ P \ n) \rrbracket \implies x \cdot_{(O_K \ P \ n)} y = x \cdot y$
apply (simp add:ring-n-pd-def Sr-def)
done

lemma (in Corps) ring-n-eSum-K-eSumTr:distinct-pds $K \ n \ P \implies$
 $(\forall j \leq m. f j \in \text{carrier } (O_K \ P \ n)) \longrightarrow \text{nsum } (O_K \ P \ n) \ f \ m = \text{nsum } K \ f \ m$
apply (induct-tac m)
apply (rule impI, simp)

apply (rule impI, simp,
subst ring-n-pd-pOp-K-pOp, assumption+,
frule-tac $n = n$ in ring-n-pd[of $n \ P$],
frule-tac Ring.ring-is-ag, drule sym, simp)
apply (rule aGroup.nsum-mem, assumption+, simp+)
done

lemma (in Corps) ring-n-eSum-K-eSum: $\llbracket \text{distinct-pds } K \ n \ P;$
 $\forall j \leq m. f j \in \text{carrier } (O_K \ P \ n) \rrbracket \implies \text{nsum } (O_K \ P \ n) \ f \ m = \text{nsum } K \ f \ m$
apply (simp add:ring-n-eSum-K-eSumTr)
done

lemma (in Corps) ideal-eSum-closed: $\llbracket \text{distinct-pds } K \ n \ P; \text{ideal } (O_K \ P \ n) \ I;$
 $\forall j \leq m. f j \in I \rrbracket \implies \text{nsum } K \ f \ m \in I$
apply (frule ring-n-pd[of $n \ P$]) **thm** Ring.ideal-nsum-closed
apply (frule-tac $n = m$ in
Ring.ideal-nsum-closed[of $(O_K \ P \ n) \ I - f$], assumption+)
apply (subst ring-n-eSum-K-eSum [THEN sym, of $n \ P \ m \ f$], assumption+,
rule allI, simp add:Ring.ideal-subset)
apply assumption
done

definition

$prime\text{-}n\text{-}pd :: [-, nat \Rightarrow ('b \Rightarrow ant) set,$
 $nat, nat] \Rightarrow 'b set$
 $(\langle 4P - - - \rangle [98, 98, 98, 99] 98) \text{ where}$
 $P_K P n j = \{x. x \in (carrier (O_K P n)) \wedge 0 < ((\nu_K (P j)) x)\}$

lemma (in Corps) $zero\text{-}in\text{-}ring\text{-}n\text{-}pd\text{-}zero\text{-}K:distinct\text{-}pds K n P \Rightarrow$

$$0_{(O_K P n)} = 0_K$$

apply (simp add:ring-n-pd-def Sr-def)

done

lemma (in Corps) $one\text{-}in\text{-}ring\text{-}n\text{-}pd\text{-}one\text{-}K:distinct\text{-}pds K n P \Rightarrow$

$$1_r(O_K P n) = 1_r$$

apply (simp add:ring-n-pd-def Sr-def)

done

lemma (in Corps) $mem\text{-}ring\text{-}n\text{-}pd\text{-}mem\text{-}K: \llbracket distinct\text{-}pds K n P; x \in carrier (O_K P n) \rrbracket$
 $\Rightarrow x \in carrier K$

apply (simp add:ring-n-pd-def Sr-def)

done

lemma (in Corps) $ring\text{-}n\text{-}tOp\text{-}K\text{-}tOp: \llbracket distinct\text{-}pds K n P; x \in carrier (O_K P n);$
 $y \in carrier (O_K P n) \rrbracket \Rightarrow x \cdot_r (O_K P n) y = x \cdot_r y$

apply (simp add:ring-n-pd-def Sr-def)

done

lemma (in Corps) $ring\text{-}n\text{-}exp\text{-}K\text{-}exp: \llbracket distinct\text{-}pds K n P; x \in carrier (O_K P n) \rrbracket$
 $\Rightarrow x^{\sim K m} = x^{\sim (O_K P n) m}$

apply (frule ring-n-pd[of n P])

apply (induct-tac m) **apply** simp

apply (simp add:one-in-ring-n-pd-one-K)

apply simp

apply (frule-tac $n = na$ in Ring.npClose[of $O_K P n x$], assumption+)

apply (simp add:ring-n-tOp-K-tOp)

done

lemma (in Corps) $prime\text{-}n\text{-}pd\text{-}prime: \llbracket distinct\text{-}pds K n P; j \leq n \rrbracket \Rightarrow$

$$prime\text{-}ideal (O_K P n) (P_K P n j)$$

apply (subst prime-ideal-def)

apply (rule conjI)

apply (simp add:ideal-def)

apply (rule conjI)

apply (rule aGroup.asubg-test)

apply (frule ring-n-pd[of n P], simp add:Ring.ring-is-ag)

apply (rule subsetI, simp add:prime-n-pd-def)

apply (subgoal-tac $0_{(O_K P n)} \in P_K P n j$)

apply blast

```

apply (simp add:zero-in-ring-n-pd-zero-K)
apply (simp add:prime-n-pd-def)
apply (simp add: ring-n-pd-def Sr-def)
apply (cut-tac field-is-ring, simp add:Ring.ring-zero)
apply (rule conjI) apply (rule allI, rule impI)
apply (cut-tac  $P = P \text{ ja}$  in representative-of-pd-valuation,
      simp add:distinct-pds-def, simp add:value-of-zero)
apply (cut-tac  $P = P \text{ j}$  in representative-of-pd-valuation,
      simp add:distinct-pds-def, simp add:value-of-zero)
apply (simp add:ant-0[THEN sym])

apply (rule ballI)+
apply (simp add:prime-n-pd-def) apply (erule conjE)+
apply (frule ring-n-pd [of  $n \text{ P}$ ], frule Ring.ring-is-ag[of  $O_K \text{ P } n$ ])
apply (frule-tac  $x = b$  in aGroup.ag-mOp-closed[of  $O_K \text{ P } n$ ], assumption+)
apply (simp add:aGroup.ag-pOp-closed)
  apply (thin-tac Ring ( $O_K \text{ P } n$ )) apply (thin-tac aGroup ( $O_K \text{ P } n$ ))
apply (simp add:ring-n-pd-def Sr-def)
apply (erule conjE)+
apply (cut-tac  $v = \nu_K (P \text{ j})$  and  $x = a$  and  $y = -_a b$  in
      amin-le-plus)
apply (rule-tac  $P = P \text{ j}$  in representative-of-pd-valuation,
      simp add:distinct-pds-def)
apply assumption+
apply (cut-tac  $P = P \text{ j}$  in representative-of-pd-valuation)
apply (simp add:distinct-pds-def)
apply (frule-tac  $x = (\nu_K (P \text{ j})) a$  and  $y = (\nu_K (P \text{ j})) (-_a b)$  in
      amin-gt[of 0])
apply (simp add:val-minus-eq)

apply (frule-tac  $y = \text{amin } ((\nu_K (P \text{ j})) a) ((\nu_K (P \text{ j})) (-_a b))$  and
       $z = (\nu_K (P \text{ j})) (a \pm -_a b)$  in aless-le-trans[of 0], assumption+)

apply (rule ballI)+
apply (frule ring-n-pd [of  $n \text{ P}$ ])
apply (frule-tac  $x = r$  and  $y = x$  in Ring.ring-tOp-closed[of  $O_K \text{ P } n$ ],
      assumption+)
apply (simp add:prime-n-pd-def)
apply (cut-tac  $P = P \text{ j}$  in representative-of-pd-valuation,
      simp add:distinct-pds-def)
apply (thin-tac Ring ( $O_K \text{ P } n$ ))
apply (simp add:prime-n-pd-def ring-n-pd-def Sr-def, (erule conjE)+,
      simp add:val-t2p)
apply (subgoal-tac  $0 \leq ((\nu_K (P \text{ j})) r)$ )
apply (simp add:aadd-pos-poss, simp)

apply (rule conjI,
      rule contrapos-pp, simp+,

```

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    simp add:prime-n-pd-def,
    (erule conjE)+, simp add: one-in-ring-n-pd-one-K,
    simp add:distinct-pds-def, (erule conjE)+,
    cut-tac representative-of-pd-valuation[of P j],
    simp add:value-of-one, simp)

apply ((rule ballI)+, rule impI)
apply (rule contrapos-pp, simp+, erule conjE,
    simp add:prime-n-pd-def, (erule conjE)+,
    simp add:ring-n-pd-def Sr-def, (erule conjE)+,
    simp add:aneg-less,
    frule-tac  $x = (\nu_K (P j)) x$  in ale-antisym[of - 0], simp,
    frule-tac  $x = (\nu_K (P j)) y$  in ale-antisym[of - 0], simp)

apply (simp add:distinct-pds-def, (erule conjE)+,
    cut-tac representative-of-pd-valuation[of P j],
    simp add:val-t2p aadd-0-l,
    simp)
done

lemma (in Corps) n-eq-val-eq-idealTr:
 $\llbracket \text{distinct-pds } K \ n \ P; x \in \text{carrier } (O_K \ P \ n); y \in \text{carrier } (O_K \ P \ n);$ 
 $\forall j \leq n. ((\nu_K (P j)) x) \leq ((\nu_K (P j)) y) \rrbracket \implies Rxa \ (O_K \ P \ n) \ y \subseteq Rxa \ (O_K \ P \ n) \ x$ 
apply (subgoal-tac  $\forall j \leq n. \text{valuation } K \ (\nu_K (P j))$ )
apply (case-tac  $x = \mathbf{0}_{(O_K \ P \ n)}$ ,
    simp add:zero-in-ring-n-pd-zero-K)
apply (simp add:value-of-zero)
apply (subgoal-tac  $y = \mathbf{0}$ , simp,
    drule-tac  $a = n$  in forall-spec, simp,
    drule-tac  $a=n$  in forall-spec, simp)
apply (cut-tac inf-ge-any[of  $(\nu_K (P \ n)) \ y$ ],
    frule ale-antisym[of  $(\nu_K (P \ n)) \ y \ \infty$ ], assumption+)
apply (rule value-inf-zero, assumption+)
apply (simp add:mem-ring-n-pd-mem-K, assumption)

apply (frule ring-n-pd[of n P])
apply (subgoal-tac  $\forall j \leq n. 0 \leq ((\nu_K (P j)) (y \cdot_r (x^{-K})))$ )
apply (subgoal-tac  $(y \cdot_r (x^{-K})) \in \text{carrier } (O_K \ P \ n)$ )
apply (cut-tac field-frac-mul[of y x],
    frule Ring.rxa-in-Rxa[of  $O_K \ P \ n \ x \ y \cdot_r (x^{-K})$ ], assumption+,
    simp add:ring-n-pd-tOp-K-tOp[THEN sym],
    frule Ring.principal-ideal[of  $O_K \ P \ n \ x$ ], assumption+)

apply (cut-tac Ring.ideal-cont-Rxa[of  $O_K \ P \ n \ (O_K \ P \ n) \ \Diamond_p \ x \ y$ ],
    assumption+,
    simp add:mem-ring-n-pd-mem-K,
    simp add:mem-ring-n-pd-mem-K,
    simp add:zero-in-ring-n-pd-zero-K)

```

apply (*frule* *Ring.rxa-in-Rxa*[*of* $O_K P n x y \cdot_r (x^{-K})$], *assumption+*,
simp *add:ring-n-pd-def* *Sr-def*,
(erule conjE)+,
cut-tac *field-is-ring*, *rule* *Ring.ring-tOp-closed*, *assumption+*,
cut-tac *invf-closed1*[*of* x], *simp*, *simp*,
simp *add:ring-n-pd-def* *Sr-def*)
apply (*cut-tac* *Ring.ring-tOp-closed*, *assumption+*,
cut-tac *field-is-ring*, *assumption+*, *simp+*,
cut-tac *invf-closed1*[*of* x], *simp*, *simp*)

apply (*rule* *allI*, *rule* *impI*, *drule-tac* $a = j$ **in** *forall-spec*, *assumption+*,
cut-tac *invf-closed1*[*of* x], *simp*, *erule conjE*)
apply (*subst* *val-t2p* [**where** $v = \nu_K P j$], *simp*,
rule *mem-ring-n-pd-mem-K*[*of* $n P y$], *assumption+*,
frule-tac $x = j$ **in** *spec*, *simp*,
simp *add:zero-in-ring-n-pd-zero-K*)
apply (*subst* *value-of-inv* [**where** $v = \nu_K P j$], *simp*,
simp *add:ring-n-pd-def* *Sr-def*, *assumption+*)
apply (*frule-tac* $x = (\nu_K (P j)) x$ **and** $y = (\nu_K (P j)) y$ **in** *ale-diff-pos*,
simp *add:diff-ant-def*,
simp *add:mem-ring-n-pd-mem-K*[*of* $n P x$] *zero-in-ring-n-pd-zero-K*)

apply (*rule* *allI*, *rule* *impI*,
simp *add:distinct-pds-def*, (*erule conjE*)+,
rule-tac $P = P j$ **in** *representative-of-pd-valuation*, *simp*)
done

lemma (**in** *Corps*) *n-eq-val-eq-ideal*: [*distinct-pds* $K n P$; $x \in \text{carrier } (O_K P n)$;
 $y \in \text{carrier } (O_K P n)$; $\forall j \leq n. ((\nu_K (P j)) x) = ((\nu_K (P j)) y)$] $\implies$
 $Rxa (O_K P n) x = Rxa (O_K P n) y$
apply (*rule* *equalityI*)
apply (*subgoal-tac* $\forall j \leq n. (\nu_K (P j)) y \leq ((\nu_K (P j)) x)$)
apply (*rule* *n-eq-val-eq-idealTr*, *assumption+*)
apply (*rule* *allI*, *rule* *impI*, *simp*)

apply (*subgoal-tac* $\forall j \leq n. (\nu_K (P j)) x \leq ((\nu_K (P j)) y)$)
apply (*rule* *n-eq-val-eq-idealTr*, *assumption+*)
apply (*rule* *allI*, *rule* *impI*)
apply *simp*
done

definition
mI-gen :: [$-$, $\text{nat} \Rightarrow ('r \Rightarrow \text{ant}) \text{ set}, \text{nat}, 'r \text{ set}] \Rightarrow 'r$ **where**
mI-gen $K P n I = (\text{SOME } x. x \in I \wedge$
 $(\forall j \leq n. (\nu_K (P j)) x = LI K (\nu_K (P j)) I))$

definition
mL :: [$-$, $\text{nat} \Rightarrow ('r \Rightarrow \text{ant}) \text{ set}, 'r \text{ set}, \text{nat}] \Rightarrow \text{int}$ **where**

```

mL K P I j = tna (LI K (νK (P j)) I)

lemma (in Corps) mI-vals-nonempty:⟦distinct-pds K n P; ideal (OK Pn) I; j ≤ n⟧
  ⇒ (νK (P j)) ' I ≠ {}
apply (frule ring-n-pd[of n P])
apply (frule Ring.ideal-zero [of OK Pn I], assumption+)

apply (simp add:image-def)
apply blast
done

lemma (in Corps) mI-vals-LB:⟦distinct-pds K n P; ideal (OK Pn) I; j ≤ n⟧ ⇒
  ((νK (P j)) ' I) ⊆ LBset (ant 0)
apply (rule subsetI)
apply (simp add:image-def, erule bexE)
  apply (frule ring-n-pd[of n P])
  apply (frule-tac h = xa in Ring.ideal-subset[of OK Pn I], assumption+)
  apply (thin-tac ideal (OK Pn) I)
  apply (thin-tac Ring (OK Pn))
  apply (simp add: ring-n-pd-def Sr-def) apply (erule conjE)+
  apply (drule-tac a = j in forall-spec, simp)

apply (simp add:LBset-def ant-0)
done

lemma (in Corps) mL-hom:⟦distinct-pds K n P; ideal (OK Pn) I;
  I ≠ {0(OK Pn)}; I ≠ carrier (OK Pn)⟧ ⇒
  ∀ j ≤ n. mL K P I j ∈ Zset
apply (rule allI, rule impI)
  apply (simp add:mL-def LI-def)
  apply (simp add:Zset-def)
done

lemma (in Corps) ex-Zleast-in-mI:⟦distinct-pds K n P; ideal (OK Pn) I; j ≤ n⟧
  ⇒ ∃ x ∈ I. (νK (P j)) x = LI K (νK (P j)) I
apply (frule-tac j = j in mI-vals-nonempty[of n P I], assumption+)
  apply (frule-tac j = j in mI-vals-LB[of n P I], assumption+)
  apply (frule-tac A = (νK (P j)) ' I and z = 0 in AMin-mem, assumption+)
  apply (simp add:LI-def)
  apply (thin-tac (νK (P j)) ' I ⊆ LBset (ant 0))
  apply (simp add:image-def, erule bexE)
  apply (drule sym)
  apply blast
done

lemma (in Corps) val-LI-pos:⟦distinct-pds K n P; ideal (OK Pn) I;
  I ≠ {0(OK Pn)}; j ≤ n⟧ ⇒ 0 ≤ LI K (νK (P j)) I
apply (frule-tac j = j in mI-vals-nonempty[of n P I], assumption+)

```


apply (frule-tac $j = j$ **in** $mI\text{-vals-LB}[of\ n\ P\ I]$, *assumption+*)
apply (frule-tac $A = (\nu_K\ (P\ j))\ 'I$ **and** $z = 0$ **in** $AMin\text{-mem}$, *assumption+*)
apply (simp add:LI-def)
apply (frule subsetD[of $(\nu_K\ (P\ j))\ 'I\ LBset\ (ant\ 0)\ AMin\ ((\nu_K\ (P\ j))\ 'I)$], *assumption+*)
apply (simp add:LBset-def ant-0)
done

lemma (**in** $Corps$) $val\text{-LI-noninf}:\llbracket distinct\text{-pds}\ K\ n\ P; ideal\ (O_K\ P\ n)\ I;$

$I \neq \{0_{(O_K\ P\ n)}\}; j \leq n \rrbracket \implies LI\ K\ (\nu_K\ (P\ j))\ I \neq \infty$

apply (frule-tac $j = j$ **in** $mI\text{-vals-nonempty}[of\ n\ P\ I]$, *assumption+*)
apply (frule-tac $j = j$ **in** $mI\text{-vals-LB}[of\ n\ P\ I]$, *assumption+*)
apply (frule-tac $A = (\nu_K\ (P\ j))\ 'I$ **and** $z = 0$ **in** $AMin$, *assumption+*)
apply (thin-tac $(\nu_K\ (P\ j))\ 'I \subseteq LBset\ (ant\ 0)$,
 $thin\text{-tac}\ (\nu_K\ (P\ j))\ 'I \neq \{\}$)
apply (frule ring-n-pd[of $n\ P$])
apply (frule Ring.ideal-zero[of $O_K\ P\ n\ I]$, *assumption+*)
apply (erule conjE, simp add:LI-def)
apply (frule singleton-sub[of $0_{O_K\ P\ n}\ I$])
apply (frule sets-not-eq[of $I\ \{0_{O_K\ P\ n}\}$],
assumption+, erule bexE)
apply (simp add:zero-in-ring-n-pd-zero-K)
apply (subgoal-tac $\exists x \in I. AMin\ ((\nu_K\ (P\ j))\ 'I) = (\nu_K\ (P\ j))\ x$,
erule bexE) **apply** simp
apply (drule-tac $x = a$ **in** $bspec$, *assumption*)
apply (thin-tac $AMin\ ((\nu_K\ (P\ j))\ 'I) = (\nu_K\ (P\ j))\ x$)

apply (frule-tac $h = a$ **in** $Ring.ideal\text{-subset}[of\ O_K\ P\ n\ I]$, *assumption+*)
apply (frule-tac $x = a$ **in** $mem\text{-ring-n-pd-mem-K}[of\ n\ P]$, *assumption+*)
apply (simp add:distinct-pds-def, (erule conjE)+)
apply (cut-tac representative-of-pd-valuation[of $P\ j$])
defer apply simp **apply** blast
apply (frule-tac $x = a$ **in** $val\text{-nonzero-z}[of\ \nu_K\ (P\ j)]$, *assumption+*,
erule exE, simp)
apply (thin-tac $\forall l \leq n. \forall m \leq n. l \neq m \longrightarrow P\ l \neq P\ m$,
 $thin\text{-tac}\ (\nu_K\ (P\ j))\ a = ant\ z$)

apply (rule contrapos-pp, simp+)
apply (cut-tac $x = ant\ z$ **in** $inf\text{-ge-any}$)
apply (frule-tac $x = ant\ z$ **in** $ale\text{-antisym}[of\ -\ \infty]$, *assumption+*)
apply simp
done

lemma (**in** $Corps$) $Zleast\text{-in-}mI\text{-pos}:\llbracket distinct\text{-pds}\ K\ n\ P; ideal\ (O_K\ P\ n)\ I;$

$I \neq \{0_{(O_K\ P\ n)}\}; j \leq n \rrbracket \implies 0 \leq mL\ K\ P\ I\ j$

apply (simp add:mL-def)
apply (frule ex-Zleast-in-mI[of $n\ P\ I\ j$], *assumption+*,

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    erule bexE, frule sym, thin-tac ( $\nu_K (P\ j)$ )  $x = LI\ K\ (\nu_K (P\ j))\ I$ )
  apply (subgoal-tac  $LI\ K\ (\nu_K (P\ j))\ I \neq \infty$ , simp)
  apply (thin-tac  $LI\ K\ (\nu_K (P\ j))\ I = (\nu_K (P\ j))\ x$ )

  apply (frule ring-n-pd[of  $n\ P$ ])
  apply (frule-tac  $h = x$  in  $Ring.ideal-subset$ [of  $O_K\ P\ n\ I$ ], assumption+)
  apply (thin-tac ideal ( $O_K\ P\ n$ )  $I$ )
  apply (thin-tac Ring ( $O_K\ P\ n$ ))
  apply (simp add: ring-n-pd-def Sr-def) apply (erule conjE)
  apply (drule-tac  $a = j$  in forall-spec, assumption)
  apply (simp add: apos-tna-pos)
  apply (rule val-LI-noninf, assumption+)
done

lemma (in Corps) Zleast-mL-I:[distinct-pds  $K\ n\ P$ ; ideal ( $O_K\ P\ n$ )  $I$ ;  $j \leq n$ ;
   $I \neq \{0_{(O_K\ P\ n)}\}; x \in I] \implies ant\ (mL\ K\ P\ I\ j) \leq ((\nu_K (P\ j))\ x)$ 
  apply (frule val-LI-pos[of  $n\ P\ I\ j$ ], assumption+)
  apply (frule apos-neq-minf[of  $LI\ K\ (\nu_K (P\ j))\ I$ ])
  apply (frule val-LI-noninf[of  $n\ P\ I\ j$ ], assumption+)
  apply (simp add: mL-def LI-def)
  apply (simp add: ant-tna)
  apply (frule Zleast-in-mL-pos[of  $n\ P\ I\ j$ ], assumption+)

  apply (frule mI-vals-nonempty[of  $n\ P\ I\ j$ ], assumption+)
  apply (frule mI-vals-LB[of  $n\ P\ I\ j$ ], assumption+)
  apply (frule AMin[of  $(\nu_K (P\ j))\ 'I\ 0$ ], assumption+)
  apply (erule conjE)
  apply (frule Zleast-in-mI-pos[of  $n\ P\ I\ j$ ], assumption+)
  apply (simp add: mL-def LI-def)
done

lemma (in Corps) Zleast-LI:[distinct-pds  $K\ n\ P$ ; ideal ( $O_K\ P\ n$ )  $I$ ;  $j \leq n$ ;
   $I \neq \{0_{(O_K\ P\ n)}\}; x \in I] \implies (LI\ K\ (\nu_K (P\ j))\ I) \leq ((\nu_K (P\ j))\ x)$ 
  apply (frule mI-vals-nonempty[of  $n\ P\ I\ j$ ], assumption+)
  apply (frule mI-vals-LB[of  $n\ P\ I\ j$ ], assumption+)
  apply (frule AMin[of  $(\nu_K (P\ j))\ 'I\ 0$ ], assumption+)
  apply (erule conjE)
  apply (simp add: LI-def)
done

lemma (in Corps) mpdiv-vals-nonequiv:distinct-pds  $K\ n\ P \implies$ 
  vals-nonequiv  $K\ n\ (\lambda j. \nu_K (P\ j))$ 
  apply (simp add: vals-nonequiv-def)
  apply (rule conjI)
  apply (simp add: valuations-def)
  apply (rule allI, rule impI)
  apply (rule representative-of-pd-valuation,
    simp add: distinct-pds-def)

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apply ((rule allI, rule impI)+, rule impI)
apply (simp add:distinct-pds-def, erule conjE)
apply (rotate-tac 4) apply (
  drule-tac a = j in forall-spec, assumption)
apply (rotate-tac -1,
  drule-tac a = l in forall-spec, assumption, simp)
apply (simp add:distinct-p-divisors)
done

```

definition

```

KbaseP :: [-, nat ⇒ ('r ⇒ ant) set, nat] ⇒
  (nat ⇒ 'r) ⇒ bool where
KbaseP K P n f ⇔ (∀ j ≤ n. f j ∈ carrier K) ∧
  (∀ j ≤ n. ∀ l ≤ n. (νK (P j)) (f l) = (δj l))

```

definition

```

Kbase :: [-, nat, nat ⇒ ('r ⇒ ant) set]
  ⇒ (nat ⇒ 'r) (⟨(3Kb_ -_)⟩ [95,95,96] 95) where
KbK n P = (SOME f. KbaseP K P n f)

```

lemma (**in** *Corps*) *KbaseTr:distinct-pds* K n P ⇒ ∃ f. *KbaseP* K P n f

```

apply (simp add: KbaseP-def)
apply (frule mpdiv-vals-nonequiv[of n P])
apply (case-tac n = 0)
apply (simp add:vals-nonequiv-def valuations-def)
apply (simp add:distinct-pds-def)
apply (frule n-val-n-val1[of P 0])
apply (frule n-val-surj[of νK (P 0)])
apply (erule bexE)
apply (subgoal-tac ((λj∈{0::nat}. x) (0::nat)) ∈ carrier K ∧
  (νK (P 0)) ((λj∈{0::nat}. x) (0::nat)) = (δ0 0))
apply blast
apply (rule conjI)
apply simp apply (simp add:Kronecker-delta-def)
apply (cut-tac Approximation1-5[of n - Suc 0 λj. νK (P j)])
apply simp
apply simp +
apply (rule allI, rule impI)
apply (rule n-val-n-val1)
apply (simp add:distinct-pds-def)
done

```

lemma (**in** *Corps*) *KbaseTr1:distinct-pds* K n P ⇒ *KbaseP* K P n (*Kb*_{K n P})

```

apply (subst Kbase-def)
apply (frule KbaseTr[of n P])
apply (erule exE)
apply (simp add:someI)
done

```

lemma (in Corps) *Kbase-hom:distinct-pds* $K\ n\ P \implies$
 $\forall j \leq n. (Kb_{K\ n\ P})\ j \in \text{carrier } K$
apply (frule *KbaseTr1*[of $n\ P$])
apply (simp add:*KbaseP-def*)
done

lemma (in Corps) *Kbase-Kronecker:distinct-pds* $K\ n\ P \implies$
 $\forall j \leq n. \forall l \leq n. (\nu_K\ (P\ j))\ ((Kb_{K\ n\ P})\ l) = \delta_j\ l$
apply (frule *KbaseTr1*[of $n\ P$])
apply (simp add:*KbaseP-def*)
done

lemma (in Corps) *Kbase-nonzero:distinct-pds* $K\ n\ P \implies$
 $\forall j \leq n. (Kb_{K\ n\ P})\ j \neq \mathbf{0}$
apply (rule *allI*, rule *impI*)
apply (frule *Kbase-Kronecker*[of $n\ P$])
apply (subgoal-tac $(\nu_K\ (P\ j))\ ((Kb_{K\ n\ P})\ j) = \delta_j\ j$)
apply (thin-tac $\forall j \leq n. (\forall l \leq n. ((\nu_K\ P\ j)\ ((Kb_{K\ n\ P})\ l)) = \delta_j\ l))$
apply (simp add:*Kronecker-delta-def*)
apply (rule *contrapos-pp*, *simp+*)
apply (cut-tac $P = P\ j$ in *representative-of-pd-valuation*)
apply (simp add:*distinct-pds-def*)
apply (simp only:*value-of-zero*, *simp only:ant-1*[*THEN sym*],
 $\text{frule } \text{sym}, \text{thin-tac } \infty = \text{ant } 1, \text{simp del:ant-1}$)
apply *simp*
done

lemma (in Corps) *Kbase-hom1:distinct-pds* $K\ n\ P \implies$
 $\forall j \leq n. (Kb_{K\ n\ P})\ j \in \text{carrier } K - \{\mathbf{0}\}$
by(simp add:*Kbase-nonzero Kbase-hom*)

definition
 $Zl-mI :: [-, \text{nat} \Rightarrow ('b \Rightarrow \text{ant})\ \text{set}, 'b\ \text{set}]$
 $\Rightarrow \text{nat} \Rightarrow 'b\ \text{where}$
 $Zl-mI\ K\ P\ I\ j = (\text{SOME } x. (x \in I \wedge ((\nu_K\ (P\ j))\ x = LI\ K\ (\nu_K\ (P\ j))\ I)))$

lemma (in Corps) *value-Zl-mI*: $\llbracket \text{distinct-pds } K\ n\ P; \text{ideal } (O_{K\ P\ n})\ I; j \leq n \rrbracket$
 $\implies (Zl-mI\ K\ P\ I\ j \in I) \wedge (\nu_K\ (P\ j))\ (Zl-mI\ K\ P\ I\ j) = LI\ K\ (\nu_K\ (P\ j))\ I$
apply (subgoal-tac $\exists x. (x \in I \wedge ((\nu_K\ (P\ j))\ x = LI\ K\ (\nu_K\ (P\ j))\ I))$)
apply (subst *Zl-mI-def*)+
apply (rule *someI2-ex*, *assumption+*)
apply (frule *ex-Zleast-in-mI*[of $n\ P\ I\ j$], *assumption+*)
apply (erule *bexE*, *blast*)
done

lemma (in Corps) *Zl-mI-nonzero*: $\llbracket \text{distinct-pds } K\ n\ P; \text{ideal } (O_{K\ P\ n})\ I;$
 $I \neq \{\mathbf{0}_{(O_{K\ P\ n})}\}; j \leq n \rrbracket \implies Zl-mI\ K\ P\ I\ j \neq \mathbf{0}$
apply (case-tac $n = 0$)

```

apply (simp add:distinct-pds-def)
apply (frule representative-of-pd-valuation[of P 0])
apply (subgoal-tac  $O_K P 0 = \text{Vr } K (\nu_K (P 0))$ )
apply (subgoal-tac  $Zl-mI K P I 0 = Ig K (\nu_K (P 0)) I$ )
apply simp apply (simp add:Ig-nonzero)
apply (simp add:Ig-def Zl-mI-def)
apply (simp add:ring-n-pd-def Vr-def)

apply (simp)
apply (frule value-Zl-mI[of n P I j], assumption+)
apply (erule conjE)
apply (rule contrapos-pp, simp+)
apply (frule distinct-pds-valuation1[of n j P], assumption+)
apply (simp add:value-of-zero)
apply (simp add:zero-in-ring-n-pd-zero-K)
apply (frule singleton-sub[of 0 I],
  frule sets-not-eq[of I {0}], assumption,
  erule bexE, simp)
apply (frule-tac  $x = a$  in Zleast-mL-I[of n P I j], assumption+)
apply (frule-tac  $x = a$  in val-nonzero-z[of  $\nu_K (P j)$ ])
apply (frule ring-n-pd[of n P])
apply (frule-tac  $h = a$  in Ring.ideal-subset[of  $O_K P_n I$ ], assumption+)
apply (simp add:mem-ring-n-pd-mem-K) apply assumption

apply (simp add:zero-in-ring-n-pd-zero-K) apply assumption
apply (frule val-LI-noninf[THEN not-sym, of n P I j], assumption+)
apply (simp add:zero-in-ring-n-pd-zero-K) apply assumption
apply simp
done

lemma (in Corps) Zl-mI-mem-K:  $\llbracket \text{distinct-pds } K n P; \text{ideal } (O_K P_n) I; l \leq n \rrbracket$ 
 $\implies (Zl-mI K P I l) \in \text{carrier } K$ 
apply (frule value-Zl-mI[of n P I l], assumption+)
apply (erule conjE)
apply (frule ring-n-pd[of n P])
apply (frule Ring.ideal-subset[of  $O_K P_n I Zl-mI K P I l$ ], assumption+)
apply (simp add:mem-ring-n-pd-mem-K[of n P Zl-mI K P I l])
done

definition
  mprod-exp ::  $[-, \text{nat} \Rightarrow \text{int}, \text{nat} \Rightarrow 'b, \text{nat}]$ 
   $\Rightarrow 'b$  where
  mprod-exp K e f n = nprod K ( $\lambda j. ((f j)_K^{(e j)})$ ) n

lemma (in Corps) mprod-expR-memTr:  $(\forall j \leq n. f j \in \text{carrier } K) \longrightarrow$ 
  mprod-expR K e f n  $\in \text{carrier } K$ 
apply (cut-tac field-is-ring)
apply (induct-tac n)
apply (rule impI, simp)

```

```

apply (simp add:mprod-expR-def)
apply (cut-tac Ring.npClose[of K f 0 e 0], assumption+)

```

```

apply (rule impI)
apply simp
apply (subst Ring.mprodR-Suc, assumption+)
apply (simp)
apply (simp)
apply (rule Ring.ring-tOp-closed[of K], assumption+)
apply (rule Ring.npClose, assumption+)
apply simp
done

```

```

lemma (in Corps) mprod-expR-mem:  $\forall j \leq n. f j \in \text{carrier } K \implies$ 
   $\text{mprod-expR } K \text{ e f } n \in \text{carrier } K$ 
apply (cut-tac field-is-ring)
apply (cut-tac Ring.mprod-expR-memTr[of K e n f])
apply simp
apply (subgoal-tac f  $\in \{j. j \leq n\} \rightarrow \text{carrier } K$ , simp+)
done

```

```

lemma (in Corps) mprod-Suc:  $\llbracket \forall j \leq (\text{Suc } n). e j \in \text{Zset};$ 
   $\forall j \leq (\text{Suc } n). f j \in (\text{carrier } K - \{0\}) \rrbracket \implies$ 
   $\text{mprod-exp } K \text{ e f } (\text{Suc } n) = (\text{mprod-exp } K \text{ e f } n) \cdot_r ((f (\text{Suc } n))_K^{(e (\text{Suc } n))})$ 
apply (simp add:mprod-exp-def)
done

```

```

lemma (in Corps) mprod-memTr:
   $(\forall j \leq n. e j \in \text{Zset}) \wedge (\forall j \leq n. f j \in ((\text{carrier } K) - \{0\})) \longrightarrow$ 
   $(\text{mprod-exp } K \text{ e f } n) \in ((\text{carrier } K) - \{0\})$ 
apply (induct-tac n)
apply (simp, rule impI, (erule conjE)+,
  simp add:mprod-exp-def, simp add:npowf-mem,
  simp add:field-potent-nonzero1)
apply (rule impI, simp, erule conjE,
  cut-tac field-is-ring, cut-tac field-is-idom,
  erule conjE, simp add:mprod-Suc)
apply (rule conjI)
apply (rule Ring.ring-tOp-closed[of K], assumption+,
  simp add:npowf-mem)
apply (rule Idomain.idom-tOp-nonzeros, assumption+,
  simp add:npowf-mem, assumption,
  simp add:field-potent-nonzero1)
done

```

```

lemma (in Corps) mprod-mem:  $\llbracket \forall j \leq n. e j \in \text{Zset}; \forall j \leq n. f j \in ((\text{carrier } K) -$ 
   $\{0\}) \rrbracket \implies (\text{mprod-exp } K \text{ e f } n) \in ((\text{carrier } K) - \{0\})$ 
apply (cut-tac mprod-memTr[of n e f]) apply simp
done

```

```

lemma (in Corps) mprod-mprodR:  $\llbracket \forall j \leq n. e\ j \in \text{Zset}; \forall j \leq n. 0 \leq (e\ j);$ 
 $\forall j \leq n. f\ j \in ((\text{carrier } K) - \{\mathbf{0}\}) \rrbracket \implies$ 
 $\text{mprod-exp } K\ e\ f\ n = \text{mprod-expR } K\ (\text{nat } o\ e)\ f\ n$ 
apply (cut-tac field-is-ring)
apply (simp add:mprod-exp-def mprod-expR-def)
apply (rule Ring.nprod-eq, assumption+)
apply (rule allI, rule impI, simp add:npowf-mem)
apply (rule allI, rule impI, rule Ring.npClose, assumption+, simp)
apply (rule allI, rule impI)
apply (simp add:npowf-def)
done

```

2.8.1 Representation of an ideal I as a product of prime ideals

```

lemma (in Corps) ring-n-mprod-mprodRTr:  $\text{distinct-pds } K\ n\ P \implies$ 
 $(\forall j \leq m. e\ j \in \text{Zset}) \wedge (\forall j \leq m. 0 \leq (e\ j)) \wedge$ 
 $(\forall j \leq m. f\ j \in \text{carrier } (O_{K\ P\ n}) - \{\mathbf{0}_{(O_{K\ P\ n})}\}) \longrightarrow$ 
 $\text{mprod-exp } K\ e\ f\ m = \text{mprod-expR } (O_{K\ P\ n})\ (\text{nat } o\ e)\ f\ m$ 
apply (frule ring-n-pd[of n P])
apply (induct-tac m)
apply (rule impI, (erule conjE)+,
  simp add:mprod-exp-def mprod-expR-def)
apply (erule conjE, simp add:npowf-def, simp add:ring-n-exp-K-exp)

apply (rule impI, (erule conjE)+, simp)
apply (subst mprod-Suc, assumption+,
  rule allI, rule impI,
  simp add:mem-ring-n-pd-mem-K,
  simp add:zero-in-ring-n-pd-zero-K)
apply (subst Ring.mprodR-Suc, assumption+,
  simp add:cmp-def,
  simp)
apply (simp add:ring-n-pd, simp add:npowf-def,
  simp add:ring-n-exp-K-exp)
apply (subst ring-n-tOp-K-tOp, assumption+,
  rule Ring.mprod-expR-mem, simp add:ring-n-pd,
  simp,
  simp)
apply (rule Ring.npClose, simp add:ring-n-pd, simp, simp)
done

```

```

lemma (in Corps) ring-n-mprod-mprodR:  $\llbracket \text{distinct-pds } K\ n\ P; \forall j \leq m. e\ j \in \text{Zset};$ 
 $\forall j \leq m. 0 \leq (e\ j); \forall j \leq m. f\ j \in \text{carrier } (O_{K\ P\ n}) - \{\mathbf{0}_{(O_{K\ P\ n})}\} \rrbracket$ 
 $\implies \text{mprod-exp } K\ e\ f\ m = \text{mprod-expR } (O_{K\ P\ n})\ (\text{nat } o\ e)\ f\ m$ 
apply (simp add:ring-n-mprod-mprodRTr)
done

```

lemma (in Corps) value-mprod-expTr:valuation K $v \implies$
 $(\forall j \leq n. e\ j \in \text{Zset}) \wedge (\forall j \leq n. f\ j \in (\text{carrier } K - \{\mathbf{0}\})) \longrightarrow$
 $v\ (\text{mprod-exp } K\ e\ f\ n) = \text{ASum } (\lambda j. (e\ j) *_a (v\ (f\ j)))\ n$
apply (induct-tac n)
apply simp
apply (rule impI, erule conjE)
apply (simp add:mprod-exp-def val-exp)

apply (rule impI, erule conjE)
apply simp
apply (subst mprod-Suc, assumption+)
apply (rule allI, rule impI, simp)
apply (subst val-t2p[of v], assumption+)
apply (cut-tac $n = n$ in mprod-mem[of - $e\ f$],
(rule allI, rule impI, simp)+, simp)
apply (simp add:npowf-mem, simp add:field-potent-nonzero1)
apply (simp add:val-exp[THEN sym, of v])
done

lemma (in Corps) value-mprod-exp:[valuation K v ; $\forall j \leq n. e\ j \in \text{Zset};$
 $\forall j \leq n. f\ j \in (\text{carrier } K - \{\mathbf{0}\})] \implies$
 $v\ (\text{mprod-exp } K\ e\ f\ n) = \text{ASum } (\lambda j. (e\ j) *_a (v\ (f\ j)))\ n$
apply (simp add:value-mprod-expTr)
done

lemma (in Corps) mgenerator0-1:[distinct-pds K (Suc n) P ;
ideal $(O_K\ P\ (\text{Suc } n))\ I$; $I \neq \{\mathbf{0}_{(O_K\ P\ (\text{Suc } n))}\}$;
 $I \neq \text{carrier } (O_K\ P\ (\text{Suc } n)); j \leq (\text{Suc } n)] \implies$
 $((\nu_K\ (P\ j))\ (\text{mprod-exp } K\ (mL\ K\ P\ I)\ (Kb_K\ (\text{Suc } n)\ P)\ (\text{Suc } n))) =$
 $((\nu_K\ (P\ j))\ (Zl-mI\ K\ P\ I\ j))$

apply (frule distinct-pds-valuation[of $j\ n\ P$], assumption+)
apply (frule mL-hom[of $\text{Suc } n\ P\ I$], assumption+)
apply (frule Kbase-hom1[of $\text{Suc } n\ P$])
apply (frule value-mprod-exp[of $\nu_K\ (P\ j)\ \text{Suc } n\ mL\ K\ P\ I$
 $Kb_K\ (\text{Suc } n)\ P$], assumption+)

apply (simp del:ASum-Suc)
apply (thin-tac $(\nu_K\ (P\ j))\ (\text{mprod-exp } K\ (mL\ K\ P\ I)\ (Kb_K\ (\text{Suc } n)\ P)\ (\text{Suc } n))$)
 $=$

$\text{ASum } (\lambda ja. (mL\ K\ P\ I\ ja) *_a (\nu_K\ (P\ j))\ ((Kb_K\ (\text{Suc } n)\ P)\ ja))\ (\text{Suc } n))$
apply (subgoal-tac ASum $(\lambda ja. (mL\ K\ P\ I\ ja) *_a$
 $((\nu_K\ (P\ j))\ ((Kb_K\ (\text{Suc } n)\ P)\ ja)))\ (\text{Suc } n) =$
 $\text{ASum } (\lambda ja. (mL\ K\ P\ I\ ja) *_a (\delta_j\ ja))\ (\text{Suc } n))$

apply (simp del:ASum-Suc)
apply (subgoal-tac $\forall h \leq (\text{Suc } n). (\lambda ja. (mL\ K\ P\ I\ ja) *_a (\delta_j\ ja))\ h \in Z_\infty)$
apply (cut-tac eSum-single[of $\text{Suc } n\ \lambda ja. (mL\ K\ P\ I\ ja) *_a (\delta_j\ ja)\ j]$)
apply simp
apply (simp add:Kronecker-delta-def asprod-n-0)


```

apply (rotate-tac -1, drule not-sym)
apply (simp add:mL-def[of K P I j])

apply (frule val-LI-noninf[of Suc n P I j], assumption+)
apply (rule not-sym, simp, simp)
apply (frule val-LI-pos[of Suc n P I j], assumption+,
  rotate-tac -2, frule not-sym, simp, simp)

apply (frule apos-neq-minf[of LI K ( $\nu_K$  (P j)) I])
apply (simp add:ant-tna)
apply (simp add:value-Zl-mI[of Suc n P I j])
apply (rule allI, rule impI)
apply (simp add:Kdelta-in-Zinf, simp)
apply (rule ballI, simp)
apply (simp add:Kronecker-delta-def, erule conjE)
apply (simp add:asprod-n-0)

apply (rule allI, rule impI)
apply (simp add:Kdelta-in-Zinf)

apply (frule Kbase-Kronecker[of Suc n P])
apply (rule ASum-eq,
  rule allI, rule impI,
  simp add:Kdelta-in-Zinf,
  rule allI, rule impI,
  simp add:Kdelta-in-Zinf)
apply (rule allI, rule impI) apply simp
done

lemma (in Corps) mgenerator0-2:[ $0 < n$ ; distinct-pds K n P; ideal ( $O_{K P n}$ ) I;
   $I \neq \{0\}_{(O_{K P n})}$ ;  $I \neq \text{carrier } (O_{K P n})$ ;  $j \leq n$ ]  $\implies$ 
 $((\nu_K (P j)) (\text{mprod-exp } K (mL K P I) (Kb_{K n P} n)) = ((\nu_K (P j)) (Zl-mI K P I j)))$ 
apply (cut-tac mgenerator0-1[of n - Suc 0 P I j])
apply simp+
done

lemma (in Corps) mgenerator1:[distinct-pds K n P; ideal ( $O_{K P n}$ ) I;
   $I \neq \{0\}_{(O_{K P n})}$ ;  $I \neq \text{carrier } (O_{K P n})$ ;  $j \leq n$ ]  $\implies$ 
 $((\nu_K (P j)) (\text{mprod-exp } K (mL K P I) (Kb_{K n P} n)) = ((\nu_K (P j)) (Zl-mI K P I j)))$ 
apply (case-tac n = 0,
  frule value-Zl-mI[of n P I j], assumption+,
  frule val-LI-noninf[of n P I j], assumption+,
  frule val-LI-pos[of n P I j], assumption+,
  frule apos-neq-minf[of LI K ( $\nu_K$  (P j)) I],
  simp add:distinct-pds-def, erule conjE)
apply (cut-tac representative-of-pd-valuation[of P j], simp+,

```

$\text{simp add:mprod-exp-def,}$
 $\text{subst val-exp[THEN sym, of } \nu_K (P \ 0) \ (Kb_{K \ 0 \ P} \ 0), \text{ assumption+,}$
 $\text{cut-tac Kbase-hom[of } 0 \ P], \text{ simp,}$
 $\text{simp add:distinct-pds-def,}$
 $\text{cut-tac Kbase-nonzero[of } 0 \ P], \text{ simp+,}$
 $\text{simp add:distinct-pds-def)}$
 $\text{apply (cut-tac Kbase-nonzero[of } 0 \ P], \text{ simp add:distinct-pds-def)}$
 $\text{apply (cut-tac Kbase-Kronecker[of } 0 \ P], \text{ simp add:distinct-pds-def)}$
 $\text{apply (simp add:Kronecker-delta-def, simp add:mL-def, simp add:ant-tna)}$
 $\text{apply (simp add:distinct-pds-def)+}$
 $\text{apply (cut-tac mgenerator0-2[of } n \ P \ I \ j], \text{ simp+})}$
 $\text{apply (simp add:distinct-pds-def) apply simp+}$
 done

lemma (in Corps) $\text{mgenerator2Tr1:} \llbracket 0 < n; j \leq n; k \leq n; \text{distinct-pds } K \ n \ P \rrbracket \implies$
 $(\nu_K (P \ j)) \ (mprod\text{-exp } K \ (\lambda l. \ \gamma_k \ l) \ (Kb_{K \ n \ P} \ n)) = (\gamma_k \ j) \ *_a \ (\delta_j \ j)$
 $\text{apply (frule distinct-pds-valuation1[of } n \ j \ P], \text{ assumption+})}$
 $\text{apply (frule K-gamma-hom[of } k \ n])}$
 $\text{apply (subgoal-tac } \forall j \leq n. \ (Kb_{K \ n \ P} \ j) \in \text{carrier } K - \{0\})}$
 $\text{apply (simp add:value-mprod-exp[of } \nu_K (P \ j) \ n \ K\text{-gamma } k \ (Kb_{K \ n \ P} \ j)])}$
 $\text{apply (subgoal-tac ASum } (\lambda ja. \ (\gamma_k \ ja) \ *_a \ (\nu_K (P \ j)) \ ((Kb_{K \ n \ P} \ ja)) \ n)$
 $\quad = \text{ASum } (\lambda ja. \ (((\gamma_k \ ja) \ *_a \ (\delta_j \ ja)))) \ n)$
 apply simp
 $\text{apply (subgoal-tac } \forall j \leq n. \ (\lambda ja. \ (\gamma_k \ ja) \ *_a \ (\delta_j \ ja)) \ j \in Z_\infty)$
 $\text{apply (cut-tac eSum-single[of } n \ \lambda ja. \ ((\gamma_k \ ja) \ *_a \ (\delta_j \ ja)) \ j], \text{ simp})}$
 $\text{apply (rule allI, rule impI, simp add:Kronecker-delta-def,}$
 $\quad \text{rule impI, simp add:asprod-n-0 Zero-in-aug-inf, assumption+})}$
 $\text{apply (rule ballI, simp)}$
 $\text{apply (simp add:K-gamma-def, rule impI, simp add:Kronecker-delta-def)}$
 $\text{apply (rule allI, rule impI)}$
 $\text{apply (simp add:Kronecker-delta-def, simp add:K-gamma-def)}$
 $\text{apply (simp add:ant-0 Zero-in-aug-inf)}$
 $\text{apply (cut-tac z-in-aug-inf[of } 1], \text{ simp add:ant-1})}$

$\text{apply (rule ASum-eq)}$
 $\text{apply (rule allI, rule impI)}$
 $\text{apply (simp add:K-gamma-def, simp add:Zero-in-aug-inf)}$
 $\text{apply (rule impI, rule value-in-aug-inf, assumption+, simp)}$
 $\text{apply (simp add:K-gamma-def Zero-in-aug-inf Kdelta-in-Zinf1)}$
 $\text{apply (rule allI, rule impI)}$
 $\text{apply (simp add:Kbase-Kronecker[of } n \ P])}$
 $\text{apply (rule Kbase-hom1, assumption+)}$
 done

lemma (in Corps) $\text{mgenerator2Tr2:} \llbracket 0 < n; j \leq n; k \leq n; \text{distinct-pds } K \ n \ P \rrbracket \implies$
 $(\nu_K (P \ j)) \ ((mprod\text{-exp } K \ (\lambda l. \ \gamma_k \ l) \ (Kb_{K \ n \ P} \ n)) K^m) = \text{ant } (m \ * \ (\gamma_k \ j))$
 $\text{apply (frule K-gamma-hom[of } k \ n])}$

apply (frule Kbase-hom1[of n P])
apply (frule mprod-mem[of n K-gamma k Kb_{K n P}], assumption+)
apply (frule distinct-pds-valuation1[of n j P], assumption+)
apply (simp, erule conjE)
apply (simp add:val-exp[THEN sym])
apply (simp add:mgenerator2Tr1)
apply (simp add:K-gamma-def Kronecker-delta-def)
apply (rule impI)
apply (simp add:asprod-def a-z-z)
done

lemma (in Corps) mgenerator2Tr3-1: $\llbracket 0 < n; j \leq n; k \leq n; j = k; \text{distinct-pds } K \ n \ P \rrbracket \implies$
 $(\nu_K (P \ j)) ((\text{mprod-exp } K \ (\lambda l. (\gamma_k \ l))) (Kb_{K \ n \ P} \ n)_{K^m}) = 0$
apply (simp add:mgenerator2Tr2) **apply** (simp add:K-gamma-def)
done

lemma (in Corps) mgenerator2Tr3-2: $\llbracket 0 < n; j \leq n; k \leq n; j \neq k; \text{distinct-pds } K \ n \ P \rrbracket \implies$
 $(\nu_K (P \ j)) ((\text{mprod-exp } K \ (\lambda l. (\gamma_k \ l))) (Kb_{K \ n \ P} \ n)_{K^m}) = \text{ant } m$
apply (simp add:mgenerator2Tr2) **apply** (simp add:K-gamma-def)
done

lemma (in Corps) mgeneratorTr4: $\llbracket 0 < n; \text{distinct-pds } K \ n \ P; \text{ideal } (O_{K \ P \ n}) \ I; I \neq \{0_{O_{K \ P \ n}}\}; I \neq \text{carrier } (O_{K \ P \ n}) \rrbracket \implies$
 $\text{mprod-exp } K \ (mL \ K \ P \ I) (Kb_{K \ n \ P} \ n) \in \text{carrier } (O_{K \ P \ n})$
apply (subst ring-n-pd-def)
apply (simp add:Sr-def)
apply (frule mL-hom[of n P I], assumption+)
apply (frule mprod-mem[of n mL K P I Kb_{K n P}])
apply (rule Kbase-hom1, assumption+)
done

apply (simp add:mprod-mem)

apply (rule allI, rule impI)
apply (simp add:mgenerator1)
apply (simp add:value-Zl-mI)
apply (simp add:val-LI-pos)
done

definition

$m\text{-zmax-pdsI-hom} :: [-, \text{nat} \Rightarrow ('b \Rightarrow \text{ant}) \text{ set}, 'b \text{ set}] \Rightarrow \text{nat} \Rightarrow \text{int}$ **where**
 $m\text{-zmax-pdsI-hom } K \ P \ I = (\lambda j. \text{tna } (A\text{Min } ((\nu_K (P \ j)) \ 'I)))$

definition

$m\text{-zmax-pdsI} :: [-, \text{nat}, \text{nat} \Rightarrow ('b \Rightarrow \text{ant}) \text{ set}, 'b \text{ set}] \Rightarrow \text{int}$ **where**
 $m\text{-zmax-pdsI } K \ n \ P \ I = (m\text{-zmax } n \ (m\text{-zmax-pdsI-hom } K \ P \ I)) + 1$

lemma (in Corps) value-Zl-mI-pos: $\llbracket 0 < n; \text{distinct-pds } K \ n \ P; \text{ideal } (O_{K \ P \ n}) \ I;$

$I \neq \{0_{(O_K P_n)}\}; I \neq \text{carrier } (O_K P_n); j \leq n; l \leq n \implies$
 $0 \leq ((\nu_K (P j)) (Zl-mI K P I l))$
apply (frule value-Zl-mI[of n P I l], assumption+)
apply (erule conjE)
apply (frule ring-n-pd[of n P])
apply (frule Ring.ideal-subset[of $O_K P_n$ I Zl-mI K P I l], assumption+)
apply (thin-tac ideal ($O_K P_n$) I)
apply (thin-tac $I \neq \{0_{O_K P_n}\}$)
apply (thin-tac $I \neq \text{carrier } (O_K P_n)$)
apply (thin-tac Ring ($O_K P_n$))
apply (simp add: ring-n-pd-def Sr-def)
done

lemma (in Corps) value-mI-genTr1: $\llbracket 0 < n; \text{distinct-pds } K n P; \text{ideal } (O_K P_n) I;$
 $I \neq \{0_{O_K P_n}\}; I \neq \text{carrier } (O_K P_n); j \leq n \rrbracket \implies$
 $(\text{mprod-exp } K (K\text{-gamma } j) (Kb_{K n P}) n)_{K^{(m\text{-zmax-pds } I K n P I)}} \in \text{carrier } K$
apply (frule K-gamma-hom[of j n])
apply (frule mprod-mem[of n K-gamma j Kb_{K n P}])
apply (rule Kbase-hom1, assumption+)
apply (rule npowf-mem)
apply simp+
done

lemma (in Corps) value-mI-genTr1-0: $\llbracket 0 < n; \text{distinct-pds } K n P;$
 $\text{ideal } (O_K P_n) I; I \neq \{0_{O_K P_n}\}; I \neq \text{carrier } (O_K P_n); j \leq n \rrbracket$
 $\implies (\text{mprod-exp } K (K\text{-gamma } j) (Kb_{K n P}) n) \in \text{carrier } K$
apply (frule K-gamma-hom[of j n])
apply (frule mprod-mem[of n K-gamma j Kb_{K n P}])
apply (rule Kbase-hom1, assumption+)
apply simp
done

lemma (in Corps) value-mI-genTr2: $\llbracket 0 < n; \text{distinct-pds } K n P; \text{ideal } (O_K P_n) I;$
 $I \neq \{0_{O_K P_n}\}; I \neq \text{carrier } (O_K P_n); j \leq n \rrbracket \implies$
 $(\text{mprod-exp } K (K\text{-gamma } j) (Kb_{K n P}) n)_{K^{(m\text{-zmax-pds } I K n P I)}} \neq 0$
apply (frule K-gamma-hom[of j n])
apply (frule mprod-mem[of n K-gamma j Kb_{K n P}])
apply (rule Kbase-hom1, assumption+) **apply** simp **apply** (erule conjE)
apply (simp add: field-potent-nonzero1)
done

lemma (in Corps) value-mI-genTr3: $\llbracket 0 < n; \text{distinct-pds } K n P; \text{ideal } (O_K P_n) I;$
 $I \neq \{0_{O_K P_n}\}; I \neq \text{carrier } (O_K P_n); j \leq n \rrbracket \implies$

$(Zl-mI\ K\ P\ I\ j) \cdot_r ((mprod-exp\ K\ (K-gamma\ j)\ (Kb_{K\ n\ P})\ n)_K^{(m-zmax-pdsI\ K\ n\ P\ I)})$
 $\in carrier\ K$
apply (*cut-tac field-is-ring*)
apply (*rule Ring.ring-tOp-closed, assumption+*)
apply (*simp add:Zl-mI-mem-K*)
apply (*simp add:value-mI-genTr1*)
done

lemma (*in Corps*) *value-mI-gen*: $\llbracket 0 < n; distinct-pds\ K\ n\ P; ideal\ (O_{K\ P\ n})\ I;$
 $I \neq \{0\}_{(O_{K\ P\ n})}; I \neq carrier\ (O_{K\ P\ n}); j \leq n \rrbracket \implies$
 $(\nu_{K\ (P\ j)} (nsum\ K\ (\lambda k. ((Zl-mI\ K\ P\ I\ k) \cdot_r ((mprod-exp\ K\ (\lambda l. (\gamma_k\ l))\ (Kb_{K\ n\ P})$
 $n)_K^{(m-zmax-pdsI\ K\ n\ P\ I)})))\ n) = LI\ K\ (\nu_{K\ (P\ j)})\ I$
apply (*cut-tac field-is-ring, frule Ring.ring-is-ag[of K]*)
apply (*case-tac j = n, simp*)
apply (*cut-tac nsum-suc[of K λk. Zl-mI K P I k ·_r*
 $mprod-exp\ K\ (K-gamma\ k)\ (Kb_{K\ n\ P})\ n_K^{m-zmax-pdsI\ K\ n\ P\ I}\ n - Suc\ 0]$,
simp,
thin-tac $\Sigma_e\ K\ (\lambda k. Zl-mI\ K\ P\ I\ k \cdot_r$
 $mprod-exp\ K\ (K-gamma\ k)\ (Kb_{K\ n\ P})\ n_K^{m-zmax-pdsI\ K\ n\ P\ I})\ n =$
 $\Sigma_e\ K\ (\lambda k. Zl-mI\ K\ P\ I\ k \cdot_r$
 $mprod-exp\ K\ (K-gamma\ k)\ (Kb_{K\ n\ P})$
 $n_K^{m-zmax-pdsI\ K\ n\ P\ I})\ (n - Suc\ 0) \pm$
 $Zl-mI\ K\ P\ I\ n \cdot_r$
 $mprod-exp\ K\ (K-gamma\ n)\ (Kb_{K\ n\ P})\ n_K^{m-zmax-pdsI\ K\ n\ P\ I})$
apply (*cut-tac distinct-pds-valuation[of n n - Suc 0 P]*)
prefer 2 apply simp
prefer 2 apply simp
apply (*subst value-less-eq1[THEN sym, of ν_K (P n)*
 $(Zl-mI\ K\ P\ I\ n) \cdot_r (mprod-exp\ K\ (K-gamma\ n)\ (Kb_{K\ n\ P})\ n_K^{m-zmax-pdsI\ K\ n\ P\ I})$
 $nsum\ K\ (\lambda k. (Zl-mI\ K\ P\ I\ k) \cdot_r (mprod-exp\ K\ (K-gamma\ k)\ (Kb_{K\ n\ P})\ n_K^{m-zmax-pdsI\ K\ n\ P\ I}))$
 $(n - Suc\ 0)]$, *assumption+*)

apply (*simp add:value-mI-genTr3*)
apply (*frule Ring.ring-is-ag[of K]*)
apply (*rule aGroup.nsum-mem[of - n - Suc 0], assumption+*)
apply (*rule allI, rule impI*)
apply (*simp add:value-mI-genTr3*)

apply (*subst val-t2p[of ν_K (P n)], assumption+*)
apply (*simp add:Zl-mI-mem-K*)
apply (*simp add:value-mI-genTr1*)

apply (*simp add:mgenerator2Tr3-1[of n n n P]*)
apply (*simp add:aadd-0-r*)
apply (*frule value-Zl-mI[of n P I n], assumption+, simp*)
apply (*erule conjE*)
apply (*frule-tac f = λk. (Zl-mI K P I k) ·_r*

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      (mprod-exp K (K-gamma k) (KbK n P) nKm-zmax-pdsI K n P I) in
      value-ge-add[of  $\nu_K(P_n)$  n - Suc 0 -
      ant (m-zmax-pdsI K n P I)])
  apply (rule allI, rule impI)
  apply (rule Ring.ring-tOp-closed, assumption+)
  apply (simp add:Zl-mI-mem-K)
  apply (simp add:value-mI-genTr1)

  apply (rule allI, rule impI) apply (simp add:cmp-def)
  apply (subst val-t2p [where v= $\nu_K P_n$ ], assumption+)
  apply (simp add:Zl-mI-mem-K)
  apply (simp add:value-mI-genTr1)

  apply (cut-tac e = K-gamma ja in mprod-mem[of n - KbK n P])
  apply (simp add:Zset-def) apply (rule Kbase-hom1, assumption+)
  apply (subst val-exp[of  $\nu_K(P_n)$ , THEN sym], assumption+)
  apply simp+

  apply (subst mgenerator2Tr1[of n n - P], assumption+, simp, simp,
  assumption+)
  apply (simp add:K-gamma-def Kronecker-delta-def)
  apply (frule-tac l = ja in value-Zl-mI-pos[of n P I n],
  assumption+, simp, simp)
  apply (simp add:Nset-preTr1)
  apply (frule-tac y = ( $\nu_K(P_n)$ ) (Zl-mI K P I ja) in
  aadd-le-mono[of 0 - ant (m-zmax-pdsI K n P I)]) apply (simp add:aadd-0-l)
  apply (subgoal-tac LI K ( $\nu_K(P_n)$ ) I < ant (m-zmax-pdsI K n P I))
  apply simp
  apply (rule aless-le-trans[of LI K ( $\nu_K(P_n)$ ) I
  ant (m-zmax-pdsI K n P I)])

  apply (simp add:m-zmax-pdsI-def)
  apply (cut-tac aless-zless[of tna (LI K ( $\nu_K(P_n)$ ) I)
  m-zmax n (m-zmax-pdsI-hom K P I) + 1])
  apply (frule val-LI-noninf[of n P I n], assumption+, simp, simp)
  apply (frule val-LI-pos[of n P I n], assumption+, simp,
  frule apos-neq-minf[of LI K ( $\nu_K(P_n)$ ) I], simp add:ant-tna)
  apply (subst m-zmax-pdsI-hom-def)
  apply (subst LI-def)
  apply (cut-tac m-zmax-gt-each[of n  $\lambda u. (tna (AMin ((\nu_K(P_u)) ' I)))$ ])
  apply simp

  apply (rule allI, rule impI)
  apply (simp add:Zset-def, simp)

  apply (subst val-t2p[of  $\nu_K(P_n)$ ], assumption+)
  apply (rule Zl-mI-mem-K, assumption+, simp)
  apply (simp add:value-mI-genTr1)

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apply (*simp add:mgenerator2Tr3-1*[of $n \ n \ n \ P \ m\text{-zmax-pdsI} \ K \ n \ P \ I$])
apply (*simp add:aadd-0-r*)
apply (*simp add:value-Zl-mI*[of $n \ P \ I \ n$])

apply (*frule aGroup.addition3*[of $K \ n - \text{Suc } 0 \ \lambda k. (Zl\text{-mI} \ K \ P \ I \ k) \cdot_r$
 $((m\text{prod-exp} \ K \ (K\text{-gamma} \ k) \ (Kb_{K \ n \ P}) \ n)_K^{(m\text{-zmax-pdsI} \ K \ n \ P \ I)}) \ j$])

apply *simp*
apply (*rule allI, rule impI*)
apply (*simp add:value-mI-genTr3*) **apply** *simp+*

apply (*thin-tac* $\Sigma_e \ K \ (\lambda k. Zl\text{-mI} \ K \ P \ I \ k \cdot_r$
 $m\text{prod-exp} \ K \ (K\text{-gamma} \ k) \ (Kb_{K \ n \ P}) \ n_K^{m\text{-zmax-pdsI} \ K \ n \ P \ I}) \ n =$
 $\Sigma_e \ K \ (cmp \ (\lambda k. Zl\text{-mI} \ K \ P \ I \ k \cdot_r$
 $m\text{prod-exp} \ K \ (K\text{-gamma} \ k) \ (Kb_{K \ n \ P}) \ n_K^{m\text{-zmax-pdsI} \ K \ n \ P \ I}) \ (\tau_j \ n)) \ n$)
apply (*cut-tac nsum-suc*[of $K \ cmp \ (\lambda k. Zl\text{-mI} \ K \ P \ I \ k \cdot_r$
 $m\text{prod-exp} \ K \ (K\text{-gamma} \ k) \ (Kb_{K \ n \ P}) \ n_K^{m\text{-zmax-pdsI} \ K \ n \ P \ I}) \ (\tau_j \ n) \ n - \text{Suc}$
 0])

apply (*simp del:nsum-suc*) **apply** (*thin-tac* $\Sigma_e \ K \ (cmp \ (\lambda k. Zl\text{-mI} \ K \ P \ I \ k \cdot_r$
 $m\text{prod-exp} \ K \ (K\text{-gamma} \ k) \ (Kb_{K \ n \ P}) \ n_K^{m\text{-zmax-pdsI} \ K \ n \ P \ I}) \ (\tau_j \ n)) \ n =$
 $\Sigma_e \ K \ (cmp \ (\lambda k. Zl\text{-mI} \ K \ P \ I \ k \cdot_r$
 $m\text{prod-exp} \ K \ (K\text{-gamma} \ k) \ (Kb_{K \ n \ P}) \ n_K^{m\text{-zmax-pdsI} \ K \ n \ P \ I}) \ (\tau_j \ n))$
 $(n - \text{Suc } 0) \pm (cmp \ (\lambda k. Zl\text{-mI} \ K \ P \ I \ k \cdot_r$
 $m\text{prod-exp} \ K \ (K\text{-gamma} \ k) \ (Kb_{K \ n \ P}) \ n_K^{m\text{-zmax-pdsI} \ K \ n \ P \ I}) \ (\tau_j \ n)) \ n$)
apply (*cut-tac distinct-pds-valuation*[of $j \ n - \text{Suc } 0 \ P$])
prefer 2 **apply** *simp* **prefer** 2 **apply** *simp*
apply (*simp add:cmp-def*)

apply (*cut-tac n-in-Nsetn*[of n])
apply (*simp add:transpos-ij-2*)
apply (*subst value-less-eq1*[*THEN sym*, of $\nu_K \ (P \ j)$
 $(Zl\text{-mI} \ K \ P \ I \ j) \cdot_r (m\text{prod-exp} \ K \ (K\text{-gamma} \ j) \ (Kb_{K \ n \ P})$
 $n_K^{m\text{-zmax-pdsI} \ K \ n \ P \ I}) \Sigma_e \ K \ (\lambda x. (Zl\text{-mI} \ K \ P \ I \ ((\tau_j \ n) \ x)) \cdot_r$
 $(m\text{prod-exp} \ K \ (K\text{-gamma} \ ((\tau_j \ n) \ x)) \ (Kb_{K \ n \ P}) \ n_K^{m\text{-zmax-pdsI} \ K \ n \ P \ I})) \ (n -$
 $\text{Suc } 0)$], *assumption+*)
apply (*simp add:value-mI-genTr3*)
apply (*rule aGroup.nsum-mem*[of $K \ n - \text{Suc } 0$], *assumption+*)
apply (*rule allI, rule impI*)
apply (*frule-tac l = ja in transpos-mem*[of $j \ n \ n$], *simp+*)
apply (*simp add:value-mI-genTr3*)

apply (*subst val-t2p*[of $\nu_K \ (P \ j)$], *assumption+*)
apply (*simp add:Zl-mI-mem-K*)
apply (*simp add:value-mI-genTr1*)

apply (*simp add:mgenerator2Tr3-1*[of $n \ j \ j \ P$])

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apply (frule value-Zl-mI[of n P I j], assumption+)
apply (erule conjE)
apply (simp add:aadd-0-r)
apply (cut-tac f =  $\lambda x. (Zl-mI\ K\ P\ I\ ((\tau_j\ n)\ x)) \cdot_r$ 
  (mprod-exp K (K-gamma (( $\tau_j\ n$ ) x)) (KbK n P) nKm-zmax-pdsI K n P I) in
  value-ge-add[of  $\nu_K\ (P\ j)$ 
    n - Suc 0 - ant (m-zmax-pdsI K n P I)], assumption+)
apply (rule allI, rule impI)
apply (frule-tac l = ja in transpos-mem[of j n n], simp+)
apply (simp add:value-mI-genTr3)
apply (rule allI, rule impI) apply (simp add:cmp-def)

apply (frule-tac l = ja in transpos-mem[of j n n], simp+)

apply (subst val-t2p [where v= $\nu_K\ P\ j$ ], assumption+)
apply (simp add:Zl-mI-mem-K)
apply (simp add:value-mI-genTr1)
apply (cut-tac k = ja in transpos-noteqTr[of n - j], simp+)
apply (subst mgenerator2Tr3-2[of n j - P], simp+)
apply (cut-tac l = ( $\tau_j\ n$ ) ja in value-Zl-mI-pos[of n P I j],
  simp+)
apply (frule-tac y = ( $\nu_K\ (P\ j)$ ) (Zl-mI K P I (( $\tau_j\ n$ ) ja)) in
  aadd-le-mono[of 0 - ant (m-zmax-pdsI K n P I)])
apply (simp add:aadd-0-l)
apply (subgoal-tac LI K ( $\nu_K\ (P\ j)$ ) I < ant (m-zmax-pdsI K n P I))
apply (rule aless-le-trans[of LI K ( $\nu_K\ (P\ j)$ ) I
  ant (m-zmax-pdsI K n P I)], assumption+)

apply (simp add:m-zmax-pdsI-def)
apply (cut-tac aless-zless[of tna (LI K ( $\nu_K\ (P\ j)$ ) I)
  m-zmax n (m-zmax-pdsI-hom K P I) + 1])
apply (frule val-LI-noninf[of n P I j], assumption+,
  frule val-LI-pos[of n P I j], assumption+,
  frule apos-neq-minf[of LI K ( $\nu_K\ (P\ j)$ ) I], simp add:ant-tna)
apply (subst m-zmax-pdsI-hom-def)
apply (subst LI-def)
apply (subgoal-tac  $\forall h \leq n. (\lambda u. (tna (AMin ((\nu_K\ (P\ u)) ' I)))) h \in Zset$ )
apply (frule m-zmax-gt-each[of n  $\lambda u. (tna (AMin ((\nu_K\ (P\ u)) ' I))))$ )
apply simp
apply (rule allI, rule impI)
apply (simp add:Zset-def)
apply (subst val-t2p[of  $\nu_K\ (P\ j)$ ], assumption+)
apply (rule Zl-mI-mem-K, assumption+)
apply (simp add:value-mI-genTr1)

apply (simp add:mgenerator2Tr3-1[of n j j P

```


$m\text{-}z\text{max}\text{-}pdsI\ K\ n\ P\ I])$

apply (*simp add:aadd-0-r*)
apply (*simp add:value-Zl-mI[of n P I j]*)
done

lemma (**in** *Corps*) *mI-gen-in-I*: $\llbracket 0 < n; \text{distinct-pds}\ K\ n\ P; \text{ideal}\ (O_K\ P\ n)\ I;$
 $I \neq \{0_{(O_K\ P\ n)}\}; I \neq \text{carrier}\ (O_K\ P\ n) \rrbracket \implies$
 $(nsum\ K\ (\lambda k. ((Zl\ mI\ K\ P\ I\ k) \cdot_r$
 $((mprod\ exp\ K\ (\lambda l. (\gamma_k\ l))\ (Kb_{K\ n\ P})\ n)_K^{(m\text{-}z\text{max}\text{-}pdsI\ K\ n\ P\ I)})))\ n) \in I$
apply (*cut-tac field-is-ring, frule ring-n-pd[of n P]*)
apply (*rule ideal-eSum-closed[of n P I n], assumption+*)
apply (*rule allI, rule impI*)
apply (*frule-tac j = j in value-Zl-mI[of n P I], assumption+*)
apply (*erule conjE*)
apply (*thin-tac* $(\nu_K\ (P\ j))\ (Zl\ mI\ K\ P\ I\ j) = LI\ K\ (\nu_K\ (P\ j))\ I$)
apply (*subgoal-tac* $(mprod\ exp\ K\ (K\text{-gamma}\ j)\ (Kb_{K\ n\ P})\ n)_K^{(m\text{-}z\text{max}\text{-}pdsI\ K\ n\ P\ I)}$
 $\in \text{carrier}\ (O_K\ P\ n))$)
apply (*frule-tac* $x = Zl\ mI\ K\ P\ I\ j$ **and**
 $r = (mprod\ exp\ K\ (K\text{-gamma}\ j)\ (Kb_{K\ n\ P})\ n)_K^{(m\text{-}z\text{max}\text{-}pdsI\ K\ n\ P\ I)}$
in *Ring.ideal-ring-multiple1*[*of* $(O_K\ P\ n)\ I$], *assumption+*)
apply (*frule-tac* $h = Zl\ mI\ K\ P\ I\ j$ **in**
 Ring.ideal-subset [*of* $O_K\ P\ n\ I$], *assumption+*)
apply (*simp add:ring-n-pd-tOp-K-tOp*[*of n P*])

apply (*subst ring-n-pd-def*) **apply** (*simp add:Sr-def*)
apply (*simp add:value-mI-genTr1*)

apply (*rule allI, rule impI*)
apply (*case-tac j = ja*)
apply (*simp add:mgenerator2Tr3-1*)

apply (*simp add:mgenerator2Tr3-2*)
apply (*simp add:m-zmax-pdsI-def*) **apply** (*simp add:m-zmax-pdsI-hom-def*)
apply (*simp only:ant-0[THEN sym]*)
apply (*simp add:aless-zless*)
apply (*subgoal-tac* $\forall l \leq n. (\lambda j. \text{tna}\ (A\text{Min}\ ((\nu_K\ (P\ j))\ 'I)))\ l \in \text{Zset}$)
apply (*frule m-zmax-gt-each*[*of n* $\lambda j. \text{tna}\ (A\text{Min}\ ((\nu_K\ (P\ j))\ 'I))$])
apply (*rotate-tac -1, drule-tac* $a = ja$ **in** *forall-spec, simp+*)
apply (*frule-tac* $j = ja$ **in** *val-LI-pos*[*of n P I*], *assumption+*)
apply (*cut-tac* $j = \text{tna}\ (LI\ K\ (\nu_K\ (P\ ja))\ I)$ **in** *ale-zle*[*of 0*])
apply (*frule-tac* $j = ja$ **in** *val-LI-noninf*[*of n P I*], *assumption+*,
 $\text{frule-tac } j = ja$ **in** *val-LI-pos*[*of n P I*], *assumption+*,
 $\text{frule-tac } a = LI\ K\ (\nu_K\ (P\ ja))\ I$ **in** *apos-neq-minf, simp add:ant-tna,*
simp add:ant-0) **apply** (*unfold LI-def*)
apply (*frule-tac* $y = \text{tna}\ (A\text{Min}\ (\nu_K\ (P\ ja))\ 'I)$ **and** $z = m\text{-}z\text{max}\ n\ (\lambda j. \text{tna}\$
 $(A\text{Min}\ (\nu_K\ (P\ j))\ 'I))$ **in** *order-trans*[*of 0*], *assumption+*)
apply (*rule-tac* $y = m\text{-}z\text{max}\ n\ (\lambda j. \text{tna}\ (A\text{Min}\ (\nu_K\ (P\ j))\ 'I))$ **and**

$z = m\text{-}zmax\ n\ (\lambda j. tna\ (AMin\ (\nu_K\ (P\ j)\ 'I))) + 1$ **in** *order-trans*[of 0],
assumption+) **apply** *simp*

apply (*rule allI*, *rule impI*) **apply** (*simp add:Zset-def*)
done

We write the element $e\Sigma\ K\ (\lambda k. (Zl\text{-}mI\ K\ P\ I\ k) \cdot_K ((mprod\text{-}exp\ K\ (K\text{-}gamma\ k)\ (Kb_{K\ n\ P})\ n)_K^{(m\text{-}zmax\text{-}pdsI\ K\ n\ P\ I)}))\ n$ as $mIg_{K\ G\ a\ i\ n\ P\ I}$

definition

$mIg :: [-, nat, nat \Rightarrow ('b \Rightarrow ant)\ set,$
 $'b\ set] \Rightarrow 'b\ (\langle 4mIg\ \dots \rangle [82,82,82,83]82)$ **where**
 $mIg_{K\ n\ P\ I} = \Sigma_e\ K\ (\lambda k. (Zl\text{-}mI\ K\ P\ I\ k) \cdot_{\tau K}$
 $((mprod\text{-}exp\ K\ (K\text{-}gamma\ k)\ (Kb_{K\ n\ P})\ n)_K^{(m\text{-}zmax\text{-}pdsI\ K\ n\ P\ I)}))\ n$

We can rewrite above two lemmas by using $mIg_{K\ G\ a\ i\ n\ P\ I}$

lemma (**in** *Corps*) *value-mI-gen1*: $\llbracket 0 < n; distinct\text{-}pds\ K\ n\ P; ideal\ (O_{K\ P\ n})\ I;$
 $I \neq \{0_{(O_{K\ P\ n})}\}; I \neq carrier\ (O_{K\ P\ n}) \rrbracket \Longrightarrow$

$\forall j \leq n. (\nu_K\ (P\ j))\ (mIg_{K\ n\ P\ I}) = LI\ K\ (\nu_K\ (P\ j))\ I$

apply (*rule allI*, *rule impI*)
apply (*simp add:mIg-def value-mI-gen*)
done

lemma (**in** *Corps*) *mI-gen-in-I1*: $\llbracket 0 < n; distinct\text{-}pds\ K\ n\ P; ideal\ (O_{K\ P\ n})\ I;$
 $I \neq \{0_{(O_{K\ P\ n})}\}; I \neq carrier\ (O_{K\ P\ n}) \rrbracket \Longrightarrow (mIg_{K\ n\ P\ I}) \in I$

apply (*simp add:mIg-def mI-gen-in-I*)
done

lemma (**in** *Corps*) *mI-principalTr*: $\llbracket 0 < n; distinct\text{-}pds\ K\ n\ P; ideal\ (O_{K\ P\ n})\ I;$
 $I \neq \{0_{(O_{K\ P\ n})}\}; I \neq carrier\ (O_{K\ P\ n}); x \in I \rrbracket \Longrightarrow$

$\forall j \leq n. ((\nu_K\ (P\ j))\ (mIg_{K\ n\ P\ I})) \leq ((\nu_K\ (P\ j))\ x)$

apply (*simp add:value-mI-gen1*)
apply (*rule allI*, *rule impI*)
apply (*rule Zleast-LI*, *assumption+*)
done

lemma (**in** *Corps*) *mI-principal*: $\llbracket 0 < n; distinct\text{-}pds\ K\ n\ P; ideal\ (O_{K\ P\ n})\ I;$
 $I \neq \{0_{(O_{K\ P\ n})}\}; I \neq carrier\ (O_{K\ P\ n}) \rrbracket \Longrightarrow$

$I = Rxa\ (O_{K\ P\ n})\ (mIg_{K\ n\ P\ I})$

apply (*frule ring-n-pd*[of $n\ P$])
apply (*rule equalityI*)
apply (*rule subsetI*)
apply (*frule-tac* $x = x$ **in** *mI-principalTr*[of $n\ P\ I$],
assumption+)
apply (*frule-tac* $y = x$ **in** *n-eq-val-eq-idealTr*[of $n\ P\ mIg_{K\ n\ P\ I}$])
apply (*frule mI-gen-in-I1*[of $n\ P\ I$], *assumption+*)
apply (*simp add:Ring.ideal-subset*)+
apply (*thin-tac* $\forall j \leq n. (\nu_K\ (P\ j))\ (mIg_{K\ n\ P\ I}) \leq (\nu_K\ (P\ j))\ x$)
apply (*frule-tac* $h = x$ **in** *Ring.ideal-subset*[of $O_{K\ P\ n}\ I$], *assumption+*)

```

apply (frule-tac  $a = x$  in Ring.a-in-principal[of  $O_K P_n$ ], assumption+)
apply (simp add:subsetD)
apply (rule Ring.ideal-cont-Rxa[of  $O_K P_n I mIg_{K_n P} I$ ], assumption+)
apply (rule mI-gen-in-I1[of  $n P I$ ], assumption+)
done

```

2.8.2 prime-n-pd

```

lemma (in Corps) prime-n-pd-principal: $\llbracket distinct-pds\ K\ n\ P; j \leq n \rrbracket \implies$ 
  ( $P_{K P_n j} = Rxa\ (O_{K P_n})\ (((Kb_{K_n P})\ j))$ )
apply (frule ring-n-pd[of  $n P$ ])
apply (frule prime-n-pd-prime[of  $n P j$ ], assumption+)
apply (simp add:prime-ideal-def, frule conjunct1)
apply (fold prime-ideal-def)
apply (thin-tac prime-ideal ( $O_{K P_n}$ ) ( $P_{K P_n j}$ ))
apply (rule equalityI)
apply (rule subsetI)
apply (frule-tac  $y = x$  in n-eq-val-eq-idealTr[of  $n P (Kb_{K_n P}) j$ ])
apply (thin-tac Ring ( $O_{K P_n}$ ), thin-tac ideal ( $O_{K P_n}$ ) ( $P_{K P_n j}$ ))
apply (simp add:ring-n-pd-def Sr-def)
apply (frule Kbase-hom[of  $n P$ ], simp)
apply (rule allI, rule impI)
apply (frule Kbase-Kronecker[of  $n P$ ])
apply (simp add:Kronecker-delta-def, rule impI)
apply (simp only:ant-0[THEN sym], simp only:ant-1[THEN sym])
apply (simp del:ant-1)
apply (simp add:prime-n-pd-def)

```

```

apply (rule allI, rule impI)
apply (frule Kbase-Kronecker[of  $n P$ ])
apply simp
apply (thin-tac  $\forall j \leq n. \forall l \leq n. (\nu_K (P j)) ((Kb_{K_n P}) l) = \delta_j l$ )
apply (case-tac  $ja = j$ , simp add:Kronecker-delta-def)
apply (thin-tac ideal ( $O_{K P_n}$ ) ( $P_{K P_n j}$ ))
apply (simp add:prime-n-pd-def, erule conjE)
apply (frule-tac  $x = x$  in mem-ring-n-pd-mem-K[of  $n P$ ],
  assumption+)
apply (case-tac  $x = \mathbf{0}_K$ )
apply (frule distinct-pds-valuation2[of  $j n P$ ], assumption+)
apply (rule gt-a0-ge-1, assumption+)

apply (simp add:Kronecker-delta-def)
apply (frule-tac  $j = ja$  in distinct-pds-valuation2[of  $- n P$ ],
  assumption+)
apply (simp add:prime-n-pd-def, erule conjE)
apply (thin-tac ideal ( $O_{K P_n}$ )  $\{x. x \in carrier\ (O_{K P_n}) \wedge 0 < (\nu_K (P j))\ x\}$ )
apply (simp add:ring-n-pd-def Sr-def)
apply (cut-tac  $h = x$  in Ring.ideal-subset[of  $O_{K P_n} P_{K P_n j}$ ])

```

apply (*frule-tac* $a = x$ **in** *Ring.a-in-principal*[*of* $O_K P n$])
apply (*simp add:Ring.ideal-subset*, *assumption*+))

apply (*rule-tac* $c = x$ **and** $A = (O_K P n) \diamond_p x$ **and** $B = (O_K P n) \diamond_p (Kb_{K n P})$
 j
in *subsetD*, *assumption*+))
apply (*simp add:Ring.a-in-principal*)
apply (*rule Ring.ideal-cont-Rxa*[*of* $O_K P n P_K P n j (Kb_{K n P}) j$], *assumption*+))
apply (*subst prime-n-pd-def*, *simp*)
apply (*frule Kbase-Kronecker*[*of* $n P$])
apply (*simp add:Kronecker-delta-def*)
apply (*simp only:ant-1*[*THEN sym*], *simp only:ant-0*[*THEN sym*])
apply (*simp del:ant-1 add:aless-zless*)
apply (*subst ring-n-pd-def*, *simp add:Sr-def*)
apply (*frule Kbase-hom*[*of* $n P$])
apply *simp*
apply (*rule allI*)
apply (*simp add:ant-0*)
apply (*rule impI*)
apply (*simp only:ant-1*[*THEN sym*], *simp only:ant-0*[*THEN sym*])
apply (*simp del:ant-1*)
done

lemma (**in** *Corps*) *ring-n-prod-primesTr*: $\llbracket 0 < n; \text{distinct-pds } K n P;$
 $\text{ideal } (O_K P n) I; I \neq \{\mathbf{0}_{O_K P n}\}; I \neq \text{carrier } (O_K P n) \rrbracket \implies$
 $\forall j \leq n. (\nu_K (P j)) (\text{mprod-exp } K (\text{mL } K P I) (Kb_{K n P}) n) =$
 $(\nu_K (P j)) (\text{mIg}_{K n P} I)$
apply (*rule allI*, *rule impI*)
apply (*simp add:mgenerator1*)
apply (*simp add:value-mI-gen1*)

apply (*simp add:value-Zl-mI*)
done

lemma (**in** *Corps*) *ring-n-prod-primesTr1*: $\llbracket 0 < n; \text{distinct-pds } K n P;$
 $\text{ideal } (O_K P n) I; I \neq \{\mathbf{0}_{O_K P n}\}; I \neq \text{carrier } (O_K P n) \rrbracket \implies$
 $I = (O_K P n) \diamond_p (\text{mprod-exp } K (\text{mL } K P I) (Kb_{K n P}) n)$
apply (*frule ring-n-pd*[*of* $n P$])
apply (*subst n-eq-val-eq-ideal*[*of* $n P \text{mprod-exp } K (\text{mL } K P I)$
 $(Kb_{K n P}) n \text{mIg}_{K n P} I$], *assumption*+))
apply (*simp add:mgeneratorTr4*)
apply (*frule mI-gen-in-I1*[*of* $n P I$], *assumption*+))
apply (*simp add:Ring.ideal-subset*)
apply (*simp add:ring-n-prod-primesTr*)
apply (*simp add:mI-principal*)
done

```

lemma (in Corps) ring-n-prod-primes:  $\llbracket 0 < n; \text{distinct-pds } K \ n \ P;$ 
  ideal  $(O_{K \ P \ n}) \ I; I \neq \{\mathbf{0}_{O_{K \ P \ n}}\}; I \neq \text{carrier } (O_{K \ P \ n});$ 
   $\forall k \leq n. J \ k = (P_{K \ P \ n} \ k) \diamond^{(O_{K \ P \ n})} (\text{nat } ((mL \ K \ P \ I) \ k)) \rrbracket \implies$ 
   $I = i\Pi_{(O_{K \ P \ n}), n} J$ 
apply (simp add: prime-n-pd-principal[of n P])
apply (subst ring-n-prod-primesTr1[of n P I], assumption+)
apply (frule ring-n-pd[of n P])
apply (frule Ring.prod-n-principal-ideal[of OK P n nat o (mL K P I) n
  KbK n P J])
apply (frule Kbase-hom[of n P])
apply (simp add: nat-def)
apply (subst ring-n-pd-def) apply (simp add: Sr-def)
apply (rule Pi-I, simp)
apply (simp add: Kbase-Kronecker[of n P])
apply (simp add: Kronecker-delta-def)
apply (simp only: ant-1[THEN sym], simp only: ant-0[THEN sym])
apply (simp del: ant-1)
apply (simp add: Kbase-hom) apply simp

apply simp
apply (frule ring-n-mprod-mprodR[of n P n mL K P I KbK n P])
apply (rule allI, rule impI, simp add: Zset-def)
apply (rule allI, rule impI)
apply (simp add: Zleast-in-mI-pos)

apply (rule allI, rule impI)
apply (subst ring-n-pd-def) apply (simp add: Sr-def)
apply (frule Kbase-hom1[of n P], simp)
apply (simp add: zero-in-ring-n-pd-zero-K)
apply (frule Kbase-Kronecker[of n P])
apply (simp add: Kronecker-delta-def)
apply (simp only: ant-1[THEN sym], simp only: ant-0[THEN sym])
apply (simp del: ant-1)
apply simp
done

end

```

```

theory Valuation3
imports Valuation2
begin

```

2.9 Completion

In this section we formalize "completion" of the ground field K

definition

```

limit :: [-, 'b  $\Rightarrow$  ant, nat  $\Rightarrow$  'b, 'b]

```

$\Rightarrow \text{bool } (\langle 4\text{lim } - - - \rangle [90, 90, 90, 91] 90) \text{ where}$
 $\text{lim}_K v f b \longleftrightarrow (\forall N. \exists M. (\forall n. M < n \longrightarrow$
 $((f n) \pm_K (-_a K b)) \in (vp K v) (Vr K v) (an N)))$

lemma *not-in-singleton-noneq*: $x \notin \{a\} \implies x \neq a$
apply *simp*
done

lemma *noneq-not-in-singleton*: $x \neq a \implies x \notin \{a\}$
apply *simp*
done

lemma *inf-neq-1*[*simp*]: $\infty \neq 1$
by (*simp only: ant-1 [THEN sym], rule z-neq-inf [THEN not-sym, of 1]*)

lemma *a1-neq-0*[*simp*]: $(1::\text{ant}) \neq 0$
by (*simp only: an-1 [THEN sym], simp only: an-0 [THEN sym],*
subst aneq-natneq [of 1 0], simp)

lemma *a1-poss*[*simp*]: $(0::\text{ant}) < 1$
by (*cut-tac zposs-aposs [of 1], simp*)

lemma *a-p1-gt*[*simp*]: $\llbracket a \neq \infty; a \neq -\infty \rrbracket \implies a < a + 1$
apply (*cut-tac aadd-poss-less [of a 1],*
simp add: aadd-commute, assumption+)
apply *simp*
done

lemma (**in** *Corps*) *vpr-pow-inter-zero*: $\text{valuation } K v \implies$
 $(\bigcap \{I. \exists n. I = (vp K v) (Vr K v) (an n)\}) = \{0\}$
apply (*frule Vr-ring [of v], frule vp-ideal [of v]*)
apply (*rule equalityI*)
defer
apply (*rule subsetI*)
apply *simp*
apply (*rule allI, rule impI, erule exE, simp*)
apply (*cut-tac n = an n in vp-apow-ideal [of v], assumption+*)
apply *simp*
apply (*cut-tac I = (vp K v) (Vr K v) (an n) in Ring.ideal-zero [of Vr K v],*
assumption+)
apply (*simp add: Vr-0-f-0*)

apply (*rule subsetI, simp*)
apply (*rule contrapos-pp, simp+*)
apply (*subgoal-tac x ∈ vp K v*)
prefer 2

```

apply (drule-tac  $x = vp\ K\ v\ (Vr\ K\ v)\ (an\ 1)$  in spec)
apply (subgoal-tac  $\exists n. vp\ K\ v\ (Vr\ K\ v)\ (an\ (Suc\ 0)) = vp\ K\ v\ (Vr\ K\ v)\ (an\ n)$ ,
  simp,
  thin-tac  $\exists n. vp\ K\ v\ (Vr\ K\ v)\ (an\ (Suc\ 0)) = vp\ K\ v\ (Vr\ K\ v)\ (an\ n)$ )
apply (simp add:r-apow-def an-def)
apply (simp only:na-1)
apply (simp only:Ring.idealpow-1-self[of Vr K v vp K v])

apply blast

apply (frule n-val-valuation[of v])

apply (frule-tac  $x = x$  in val-nonzero-z[of n-val K v],
  frule-tac  $h = x$  in Ring.ideal-subset[of Vr K v vp K v],
  assumption+,
  simp add:Vr-mem-f-mem, assumption+) apply (
  frule-tac  $h = x$  in Ring.ideal-subset[of Vr K v vp K v],
  simp add:Vr-mem-f-mem, assumption+)
apply (cut-tac  $x = x$  in val-pos-mem-Vr[of v], assumption) apply(
  simp add:Vr-mem-f-mem, simp)
apply (frule-tac  $x = x$  in val-pos-n-val-pos[of v],
  simp add:Vr-mem-f-mem, simp)
apply (cut-tac  $x = n\text{-val}\ K\ v\ x$  and  $y = 1$  in aadd-pos-poss, assumption+,
  simp) apply (frule-tac  $y = n\text{-val}\ K\ v\ x + 1$  in aless-imp-le[of 0])
apply (cut-tac  $x1 = x$  and  $n1 = (n\text{-val}\ K\ v\ x) + 1$  in n-val-n-pow[THEN sym,
  of v], assumption+)
apply (drule-tac  $a = vp\ K\ v\ (Vr\ K\ v)\ (n\text{-val}\ K\ v\ x + 1)$  in forall-spec)
apply (erule exE, simp)
apply (simp only:ant-1[THEN sym] a-zpz,
  cut-tac  $z = z + 1$  in z-neq-inf)
apply (subst an-na[THEN sym], assumption+, blast)
apply simp
apply (cut-tac  $a = n\text{-val}\ K\ v\ x$  in a-p1-gt)
apply (erule exE, simp only:ant-1[THEN sym], simp only:a-zpz z-neq-inf)
apply (cut-tac  $i = z + 1$  and  $j = z$  in ale-zle, simp)
apply (erule exE, simp add:z-neq-minf)
apply (cut-tac  $y1 = n\text{-val}\ K\ v\ x$  and  $x1 = n\text{-val}\ K\ v\ x + 1$  in
  aneg-le[THEN sym], simp)
done

lemma (in Corps) limit-diff-n-val:[ $b \in carrier\ K; \forall j. f\ j \in carrier\ K;$ 
  valuation  $K\ v$ ]  $\implies (lim_{K\ v}\ f\ b) = (\forall N. \exists M. \forall n. M < n \longrightarrow$ 
   $(an\ N) \leq (n\text{-val}\ K\ v\ ((f\ n) \pm (-_a\ b))))$ 
apply (rule iffI)
apply (rule allI)
apply (simp add:limit-def) apply (rotate-tac -1)
apply (drule-tac  $x = N$  in spec)
apply (erule exE)
apply (subgoal-tac  $\forall n > M. (an\ N) \leq (n\text{-val}\ K\ v\ (f\ n \pm (-_a\ b))))$ 

```

```

apply blast
apply (rule allI, rule impI) apply (rotate-tac -2)
apply (drule-tac  $x = n$  in spec, simp)
apply (rule n-value-x-1 [of v], assumption+,
      simp add:an-nat-pos, assumption)

apply (simp add:limit-def)
apply (rule allI, rotate-tac -1, drule-tac  $x = N$  in spec)

apply (erule exE)
apply (subgoal-tac  $\forall n > M. f\ n \pm (-_a\ b) \in vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$ )
apply blast
apply (rule allI, rule impI)
apply (rotate-tac -2, drule-tac  $x = n$  in spec, simp)
apply (cut-tac field-is-ring, frule Ring.ring-is-ag [of K])
apply (rule-tac  $x = f\ n \pm -_a\ b$  and  $n = an\ N$  in n-value-x-2 [of v],
      assumption+)
apply (subst val-pos-mem-Vr [THEN sym, of v], assumption+)
apply (drule-tac  $x = n$  in spec)
apply (rule aGroup.ag-pOp-closed [of K], assumption+)
apply (simp add:aGroup.ag-mOp-closed)
apply (subst val-pos-n-val-pos [of v], assumption+)
apply (rule aGroup.ag-pOp-closed, assumption+) apply simp
apply (simp add:aGroup.ag-mOp-closed)
apply (rule-tac  $j = an\ N$  and  $k = n\text{-val}\ K\ v\ (f\ n \pm -_a\ b)$  in
      ale-trans [of 0], simp, assumption+)
apply simp
done

lemma (in Corps) an-na-Lv:valuation  $K\ v \implies an\ (na\ (Lv\ K\ v)) = Lv\ K\ v$ 
apply (frule Lv-pos [of v])
apply (frule aless-imp-le [of 0 Lv K v])
apply (frule apos-neq-minf [of Lv K v])
apply (frule Lv-z [of v], erule exE)
apply (rule an-na)
apply (rule contraposp-pp, simp+)
done

lemma (in Corps) limit-diff-val:  $\llbracket b \in carrier\ K; \forall j. f\ j \in carrier\ K; \text{valuation}\ K\ v \rrbracket \implies$ 
   $(\lim_{K\ v} f\ b) = (\forall N. \exists M. \forall n. M < n \longrightarrow$ 
     $(an\ N) \leq (v\ ((f\ n) \pm (-_a\ b))))$ 
apply (cut-tac field-is-ring, frule Ring.ring-is-ag [of K],
      simp add:limit-diff-n-val [of b f v])
apply (rule iffI)
apply (rule allI,
      rotate-tac -1, drule-tac  $x = N$  in spec, erule exE)
apply (subgoal-tac  $\forall n > M. an\ N \leq v\ (f\ n \pm -_a\ b)$ , blast)
apply (rule allI, rule impI)
apply (drule-tac  $x = n$  in spec,

```



```

    drule-tac x = n in spec, simp)
  apply (frule aGroup.ag-mOp-closed[of K b], assumption+,
    frule-tac x = f n and y = -a b in aGroup.ag-pOp-closed, assumption+)
  apply (frule-tac x = f n ± -a b in n-val-le-val[of v],
    assumption+)
  apply (cut-tac n = N in an-nat-pos)
  apply (frule-tac j = an N and k = n-val K v ( f n ± -a b) in
    ale-trans[of 0], assumption+)
  apply (subst val-pos-n-val-pos, assumption+)
  apply (rule-tac i = an N and j = n-val K v ( f n ± -a b) and
    k = v ( f n ± -a b) in ale-trans, assumption+)
  apply (rule allI)
  apply (rotate-tac 3, drule-tac x = N * (na (Lv K v)) in spec)
  apply (erule exE)
  apply (subgoal-tac ∀ n. M < n ⟶ (an N) ≤ n-val K v (f n ± -a b), blast)
  apply (rule allI, rule impI)
  apply (rotate-tac -2, drule-tac x = n in spec, simp)

  apply (drule-tac x = n in spec)
  apply (frule aGroup.ag-mOp-closed[of K b], assumption+,
    frule-tac x = f n and y = -a b in aGroup.ag-pOp-closed, assumption+)
  apply (cut-tac n = N * na (Lv K v) in an-nat-pos)
  apply (frule-tac i = 0 and j = an (N * na (Lv K v)) and
    k = v ( f n ± -a b) in ale-trans, assumption+)
  apply (simp add:amult-an-an, simp add:an-na-Lv)
  apply (frule Lv-pos[of v])
  apply (frule-tac x1 = f n ± -a b in n-val[THEN sym, of v],
    assumption+, simp)
  apply (frule Lv-z[of v], erule exE, simp)
  apply (simp add:amult-pos-mono-r)
done

```

uniqueness of the limit is derived from *vp-pow-inter-zero*

```

lemma (in Corps) limit-unique: [b ∈ carrier K; ∀ j. f j ∈ carrier K;
  valuation K v; c ∈ carrier K; limK v f b; limK v f c] ⟹ b = c
  apply (rule contrapos-pp, simp+, simp add:limit-def,
    cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    simp add:aGroup.ag-neq-diffnonzero[of K b c],
    frule vpr-pow-inter-zero[THEN sym, of v],
    frule noneq-not-in-singleton[of b ± -a c 0])
  apply simp
  apply (rotate-tac -1, erule exE, erule conjE)
  apply (erule exE, simp, thin-tac x = (vp K v) (Vr K v) (an n))
  apply (rotate-tac 3, drule-tac x = n in spec,
    erule exE,
    drule-tac x = n in spec,
    erule exE)
  apply (rename-tac x N M1 M2)
  apply (subgoal-tac M1 < Suc (max M1 M2),

```

```

    subgoal-tac M2 < Suc (max M1 M2),
    drule-tac x = Suc (max M1 M2) in spec,
    drule-tac x = Suc (max M1 M2) in spec,
    drule-tac x = Suc (max M1 M2) in spec)
apply simp

apply (frule-tac n = an N in vp-apow-ideal[of v],
    frule-tac x = f (Suc (max M1 M2))  $\pm$   $-_a$  b and N = an N in
    mem-vp-apow-mem-Vr[of v], simp,
    frule Vr-ring[of v], simp, simp)

apply (frule Vr-ring[of v],
    frule-tac I = vp K v (Vr K v) (an N) and x = f (Suc (max M1 M2))  $\pm$   $-_a$  b
    in Ring.ideal-inv1-closed[of Vr K v], assumption+)

apply (frule-tac I = vp K v (Vr K v) (an N) and h = f (Suc (max M1 M2))  $\pm$   $-_a$ 
b
in Ring.ideal-subset[of Vr K v], assumption+)
apply (simp add: Vr-mOp-f-mOp,
    cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    frule aGroup.ag-mOp-closed[of K b], assumption+,
    simp add:aGroup.ag-p-inv, simp add:aGroup.ag-inv-inv)

apply (frule-tac I = vp K v (Vr K v) (an N) and x =
 $-_a$  (f (Suc (max M1 M2)))  $\pm$  b and y = f (Suc (max M1 M2))  $\pm$  ( $-_a$  c) in
Ring.ideal-pOp-closed[of Vr K v], assumption+)
apply (frule-tac x = f (Suc (max M1 M2))  $\pm$   $-_a$  c and N = an N in
    mem-vp-apow-mem-Vr[of v], simp, assumption,
    frule-tac x =  $-_a$  (f (Suc (max M1 M2)))  $\pm$  b and N = an N in
    mem-vp-apow-mem-Vr[of v], simp,
    simp add: Vr-pOp-f-pOp, simp add: Vr-pOp-f-pOp,
    frule aGroup.ag-mOp-closed[of K c], assumption+,
    frule-tac x = f (Suc (max M1 M2)) in aGroup.ag-mOp-closed[of K],
    assumption+,
    frule-tac x = f (Suc (max M1 M2)) and y =  $-_a$  c in
    aGroup.ag-pOp-closed[of K], assumption+)

apply (simp add:aGroup.ag-pOp-assoc[of K - b -],
    simp add:aGroup.ag-pOp-assoc[THEN sym, of K b -  $-_a$  c],
    simp add:aGroup.ag-pOp-commute[of K b],
    simp add:aGroup.ag-pOp-assoc[of K - b  $-_a$  c],
    frule aGroup.ag-pOp-closed[of K b  $-_a$  c], assumption+,
    simp add:aGroup.ag-pOp-assoc[THEN sym, of K - - b  $\pm$   $-_a$  c],
    simp add:aGroup.ag-l-inv1, simp add:aGroup.ag-l-zero)
apply (simp add:aGroup.ag-pOp-commute[of K - b])
apply arith apply arith
done

```

```

lemma (in Corps) limit-n-val: $\llbracket b \in \text{carrier } K; b \neq \mathbf{0}; \text{valuation } K \ v;$ 
 $\forall j. f \ j \in \text{carrier } K; \lim_{K \ v} f \ b \rrbracket \implies$ 
 $\exists N. (\forall m. N < m \longrightarrow (n\text{-val } K \ v) (f \ m) = (n\text{-val } K \ v) \ b)$ 
apply (simp add:limit-def)
apply (frule n-val-valuation[of v])
apply (frule val-nonzero-z[of n-val K v b], assumption+, erule exE,
rename-tac L)
apply (rotate-tac -3, drule-tac x = nat (abs L + 1) in spec)
apply (erule exE, rename-tac M)

apply (subgoal-tac  $\forall n. M < n \longrightarrow n\text{-val } K \ v \ (f \ n) = n\text{-val } K \ v \ b$ , blast)
apply (rule allI, rule impI)
apply (rotate-tac -2,
drule-tac x = n in spec,
simp)
apply (frule-tac x = f n  $\pm$   $-_a$  b and n = an (nat (|L| + 1)) in
n-value-x-1[of v],
thin-tac f n  $\pm$   $-_a$  b  $\in$  vp K v (Vr K v) (an (nat (|L| + 1))))
apply (simp add:an-def)
apply (simp add:zpos-apos [symmetric])
apply assumption

apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (frule aGroup.ag-mOp-closed[of K b], assumption+)

apply (drule-tac x = n in spec,
frule-tac x = f n in aGroup.ag-pOp-closed[of K -  $-_a$  b],
assumption+,
frule-tac x = b and y = (f n)  $\pm$  ( $-_a$  b) in value-less-eq[of
n-val K v], assumption+)
apply simp
apply (rule-tac x = ant L and y = an (nat (|L| + 1)) and
z = n-val K v ( f n  $\pm$   $-_a$  b) in aless-le-trans)
apply (subst an-def)
apply (subst aless-zless) apply arith apply assumption+
apply (simp add:aGroup.ag-pOp-commute[of K b])
apply (simp add:aGroup.ag-pOp-assoc) apply (simp add:aGroup.ag-l-inv1)
apply (simp add:aGroup.ag-r-zero)
done

lemma (in Corps) limit-val: $\llbracket b \in \text{carrier } K; b \neq \mathbf{0}; \forall j. f \ j \in \text{carrier } K;$ 
 $\text{valuation } K \ v; \lim_{K \ v} f \ b \rrbracket \implies \exists N. (\forall n. N < n \longrightarrow v \ (f \ n) = v \ b)$ 
apply (frule limit-n-val[of b v f], assumption+)
apply (erule exE)
apply (subgoal-tac  $\forall m. N < m \longrightarrow v \ (f \ m) = v \ b$ )
apply blast

```

```

apply (rule allI, rule impI)
apply (drule-tac x = m in spec)
apply (drule-tac x = m in spec)
apply simp
apply (frule Lv-pos[of v])
apply (simp add:n-val[THEN sym, of v])
done

```

```

lemma (in Corps) limit-val-infinity:⟦valuation K v; ∀j. f j ∈ carrier K;
  limK v f 0⟧ ⇒ ∀N. (∃M. (∀m. M < m → (an N) ≤ (n-val K v (f m))))
apply (simp add:limit-def)
apply (rule allI)
apply (drule-tac x = N in spec,
  erule exE)

```

```

apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
  simp only:aGroup.ag-inv-zero[of K], simp only:aGroup.ag-r-zero)
apply (subgoal-tac ∀n. M < n → an N ≤ n-val K v (f n), blast)

```

```

apply (rule allI, rule impI)
apply (drule-tac x = n in spec,
  drule-tac x = n in spec, simp)
apply (simp add:n-value-x-1)
done

```

```

lemma (in Corps) not-limit-zero:⟦valuation K v; ∀j. f j ∈ carrier K;
  ¬ (limK v f 0)⟧ ⇒ ∃N. (∀M. (∃m. (M < m) ∧
    ((n-val K v (f m)) < (an N))))
apply (simp add:limit-def)
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (simp add:aGroup.ag-inv-zero aGroup.ag-r-zero)
apply (erule exE)
apply (case-tac N = 0, simp add:r-apow-def)
apply (subgoal-tac ∀M. ∃n. M < n ∧ n-val K v (f n) < (an 0)) apply blast
apply (rule allI,
  rotate-tac 4, frule-tac x = M in spec,
  erule exE, erule conjE)
apply (frule-tac x1 = f n in val-pos-mem-Vr[THEN sym, of v]) apply simp
apply simp
apply (frule-tac x = f n in val-pos-n-val-pos[of v])
apply simp apply simp apply (thin-tac ¬ 0 ≤ v (f n))
apply (simp add:aneg-le) apply blast

```

```

apply (simp)
apply (subgoal-tac ∀n. ((f n) ∈ vp K v (Vr K v) (an N)) =
  ((an N) ≤ n-val K v (f n)))
apply simp
apply (thin-tac ∀n. (f n ∈ vp K v (Vr K v) (an N)) = (an N ≤ n-val K v (f n)))
apply (simp add:aneg-le) apply blast

```

apply (*rule allI*)
apply (*thin-tac* $\forall M. \exists n. M < n \wedge f n \notin vp K v (Vr K v) (an N)$)
apply (*rule iffI*)
apply (*frule-tac* $x1 = f n$ **and** $n1 = an N$ **in** *n-val-n-pow*[*THEN sym, of v*])
apply (*rule-tac* $N = an N$ **and** $x = f n$ **in** *mem-vp-apow-mem-Vr*[*of v*],
assumption+, *simp*, *assumption*, *simp*, *simp*)
apply (*frule-tac* $x = f n$ **and** $n = an N$ **in** *n-val-n-pow*[*of v*])
apply (*frule-tac* $x1 = f n$ **in** *val-pos-mem-Vr*[*THEN sym, of v*])
apply *simp*
apply (*cut-tac* $n = N$ **in** *an-nat-pos*)
apply (*frule-tac* $j = an N$ **and** $k = n-val K v (f n)$ **in** *ale-trans*[*of 0*],
assumption+)
apply (*frule-tac* $x1 = f n$ **in** *val-pos-n-val-pos*[*THEN sym, of v*])
apply *simp+*
done

lemma (*in Corps*) *limit-p*: $\llbracket valuation K v; \forall j. f j \in carrier K;$
 $\forall j. g j \in carrier K; b \in carrier K; c \in carrier K; \lim_{K v} f b; \lim_{K v} g c \rrbracket$
 $\implies \lim_{K v} (\lambda j. (f j) \pm (g j)) (b \pm c)$
apply (*simp add:limit-def*)
apply (*rule allI*) **apply** (*rotate-tac* 3,
drule-tac $x = N$ **in** *spec*,
drule-tac $x = N$ **in** *spec*,
(erule exE)+)
apply (*frule-tac* $x = M$ **and** $y = Ma$ **in** *two-inequalities*, *simp*,
subgoal-tac $\forall n > (max M Ma). (f n) \pm (g n) \pm -_a (b \pm c)$
 $\in (vp K v)(Vr K v) (an N)$)
apply (*thin-tac* $\forall n. Ma < n \longrightarrow$
 $g n \pm -_a c \in (vp K v)(Vr K v) (an N),$
thin-tac $\forall n. M < n \longrightarrow$
 $f n \pm -_a b \in (vp K v)(Vr K v) (an N),$ *blast*)
apply (*frule Vr-ring*[*of v*],
frule-tac $n = an N$ **in** *vp-apow-ideal*[*of v*])
apply *simp*
apply (*rule allI*, *rule impI*)
apply (*thin-tac* $\forall n > M. f n \pm -_a b \in vp K v (Vr K v) (an N),$
thin-tac $\forall n > Ma. g n \pm -_a c \in vp K v (Vr K v) (an N),$
frule-tac $I = vp K v (Vr K v) (an N)$ **and** $x = f n \pm -_a b$
and $y = g n \pm -_a c$ **in** *Ring.ideal-pOp-closed*[*of Vr K v*],
assumption+, *simp*, *simp*)
apply (*frule-tac* $I = vp K v (Vr K v) (an N)$ **and** $h = f n \pm -_a b$
in *Ring.ideal-subset*[*of Vr K v*], *assumption+*, *simp*,
frule-tac $I = vp K v (Vr K v) (an N)$ **and** $h = g n \pm -_a c$ **in**
Ring.ideal-subset[*of Vr K v*], *assumption+*, *simp*)
apply (*simp add:Vr-pOp-f-pOp*)
apply (*thin-tac* $f n \pm -_a b \in carrier (Vr K v),$
thin-tac $g n \pm -_a c \in carrier (Vr K v)$)

apply (*cut-tac field-is-ring*, *frule Ring.ring-is-ag*[*of K*],
frule aGroup.ag-mOp-closed[*of K b*], *assumption+*,
frule aGroup.ag-mOp-closed[*of K c*], *assumption+*,
frule-tac a = f n and b = -_a b and c = g n and d = -_a c in
aGroup.ag-add4-rel[*of K*], *simp+*)
apply (*simp add:aGroup.ag-p-inv*[*THEN sym*])
done

lemma (*in Corps*) *Abs-ant-abs*[*simp*]:*Abs* (*ant z*) = *ant* (*abs z*)
apply (*simp add:Abs-def*, *simp add:aminus*)
apply (*simp only:ant-0*[*THEN sym*], *simp only:aless-zless*)
apply (*simp add:zabs-def*)
done

lemma (*in Corps*) *limit-t-nonzero*: $\llbracket$ *valuation K v*; $\forall (j::\text{nat}). (f j) \in \text{carrier } K$;
 $\forall (j::\text{nat}). g j \in \text{carrier } K$; $b \in \text{carrier } K$; $c \in \text{carrier } K$; $b \neq \mathbf{0}$; $c \neq \mathbf{0}$;
 $\lim_K v f b$; $\lim_K v g c \rrbracket \implies \lim_K v (\lambda j. (f j) \cdot_r (g j)) (b \cdot_r c)$
apply (*cut-tac field-is-ring*, *simp add:limit-def*, *rule allI*)
apply (*subgoal-tac* $\forall j. (f j) \cdot_r (g j) \pm -_a (b \cdot_r c) =$
 $((f j) \cdot_r (g j) \pm (f j) \cdot_r (-_a c)) \pm ((f j) \cdot_r c \pm -_a (b \cdot_r c))$, *simp*,
thin-tac $\forall j. f j \cdot_r g j \pm -_a b \cdot_r c =$
 $f j \cdot_r g j \pm f j \cdot_r (-_a c) \pm (f j \cdot_r c \pm -_a b \cdot_r c)$)
apply (*frule limit-n-val*[*of b v f*], *assumption+*,
simp add:limit-def)
apply (*frule n-val-valuation*[*of v*])
apply (*frule val-nonzero-z*[*of n-val K v b*], *assumption+*,
rotate-tac -1, *erule exE*,
frule val-nonzero-z[*of n-val K v c*], *assumption+*,
rotate-tac -1, *erule exE*, *rename-tac N bz cz*)
apply (*rotate-tac 5*,
drule-tac $x = N + \text{nat } (abs cz)$ **in spec**,
drule-tac $x = N + \text{nat } (abs bz)$ **in spec**)
apply (*erule exE*)
apply (*rename-tac N bz cz M M1 M2*)

apply (*subgoal-tac* $\forall n. (\max (\max M1 M2) M) < n \longrightarrow$
 $(f n) \cdot_r (g n) \pm (f n) \cdot_r (-_a c) \pm ((f n) \cdot_r c \pm (-_a (b \cdot_r c)))$
 $\in vp K v (Vr K v) (an N)$, *blast*)
apply (*rule allI*, *rule impI*) **apply** (*simp*, (*erule conjE*)
apply (*rotate-tac 11*, *drule-tac* $x = n$ **in spec**,
drule-tac $x = n$ **in spec**, *simp*,
drule-tac $x = n$ **in spec**, *simp*)
apply (*frule-tac* $b = g n \pm -_a c$ **and** $n = an N$ **and** $x = f n$ **in**
convergenceTr1[*of v*])
apply *simp* **apply** *simp* **apply** (*simp add:an-def a-zpz*[*THEN sym*]) **apply** *simp*
apply (*frule-tac* $b = f n \pm -_a b$ **and** $n = an N$ **in** *convergenceTr1*[*of*
 $v c$], *assumption+*, *simp*) **apply** (*simp add:an-def*)
apply (*simp add:a-zpz*[*THEN sym*]) **apply** *simp*

```

apply (drule-tac  $x = n$  in spec,
        drule-tac  $x = n$  in spec)
apply (simp add:Ring.ring-inv1-1[of  $K$   $b$   $c$ ],
        cut-tac Ring.ring-is-ag, frule aGroup.ag-mOp-closed[of  $K$   $c$ ],
        assumption+,
        simp add:Ring.ring-distrib1[THEN sym],
        frule aGroup.ag-mOp-closed[of  $K$   $b$ ], assumption+,
        simp add:Ring.ring-distrib2[THEN sym],
        subst Ring.ring-tOp-commute[of  $K$  -  $c$ ], assumption+,
        rule aGroup.ag-pOp-closed, assumption+)
apply (cut-tac  $n = N$  in an-nat-pos)
apply (frule-tac  $n = an$   $N$  in vp-apow-ideal[of  $v$ ], assumption+)
apply (frule Vr-ring[of  $v$ ])

apply (frule-tac  $x = (f\ n) \cdot_r (g\ n \pm -_a\ c)$  and  $y = c \cdot_r (f\ n \pm -_a\ b)$ 
        and  $I = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$  in Ring.ideal-pOp-closed[of  $Vr\ K\ v$ ],
        assumption+)
apply (frule-tac  $R = Vr\ K\ v$  and  $I = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$  and
         $h = (f\ n) \cdot_r (g\ n \pm -_a\ c)$  in Ring.ideal-subset, assumption+,
        frule-tac  $R = Vr\ K\ v$  and  $I = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$  and
         $h = c \cdot_r (f\ n \pm -_a\ b)$  in Ring.ideal-subset, assumption+)
apply (simp add:Vr-pOp-f-pOp, assumption+)
apply (rule allI)
apply (thin-tac  $\forall N. \exists M. \forall n > M. f\ n \pm -_a\ b \in vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$ ,
        thin-tac  $\forall N. \exists M. \forall n > M. g\ n \pm -_a\ c \in vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$ )
apply (drule-tac  $x = j$  in spec,
        drule-tac  $x = j$  in spec,
        cut-tac field-is-ring, frule Ring.ring-is-ag[of  $K$ ],
        frule-tac  $x = f\ j$  and  $y = g\ j$  in Ring.ring-tOp-closed, assumption+,
        frule-tac  $x = b$  and  $y = c$  in Ring.ring-tOp-closed, assumption+,
        frule-tac  $x = f\ j$  and  $y = c$  in Ring.ring-tOp-closed, assumption+,
        frule-tac  $x = c$  in aGroup.ag-mOp-closed[of  $K$ ], assumption+,
        frule-tac  $x = f\ j$  and  $y = -_a\ c$  in Ring.ring-tOp-closed, assumption+,
        frule-tac  $x = b \cdot_r c$  in aGroup.ag-mOp-closed[of  $K$ ], assumption+)
apply (subst aGroup.pOp-assocTr41[THEN sym, of  $K$ ], assumption+,
        subst aGroup.pOp-assocTr42[of  $K$ ], assumption+,
        subst Ring.ring-distrib1[THEN sym, of  $K$ ], assumption+)
apply (simp add:aGroup.ag-l-inv1, simp add:Ring.ring-times-x-0,
        simp add:aGroup.ag-r-zero)
done

lemma an-npn[simp]:  $an\ (n + m) = an\ n + an\ m$ 
by (simp add:an-def a-zpz)

lemma Abs-noninf:  $a \neq -\infty \wedge a \neq \infty \implies Abs\ a \neq \infty$ 
by (cut-tac mem-ant[of  $a$ ], simp, erule exE, simp add:Abs-def,
        simp add:aminus)

```

lemma (in Corps) limit-t-zero: $\llbracket c \in \text{carrier } K; \text{valuation } K v;$
 $\forall (j::\text{nat}). f j \in \text{carrier } K; \forall (j::\text{nat}). g j \in \text{carrier } K;$
 $\lim_{K v} f \mathbf{0}; \lim_{K v} g c \rrbracket \implies \lim_{K v} (\lambda j. (f j) \cdot_r (g j)) \mathbf{0}$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
simp add:limit-def, simp add:aGroup.ag-inv-zero, simp add:aGroup.ag-r-zero)
apply (subgoal-tac $\forall j. (f j) \cdot_r (g j) \in \text{carrier } K,$
simp add:aGroup.ag-r-zero)
prefer 2 **apply** (rule allI, simp add:Ring.ring-tOp-closed)
apply (case-tac $c = \mathbf{0}_K$)
apply (simp add:aGroup.ag-inv-zero, simp add:aGroup.ag-r-zero)
apply (rule allI,
rotate-tac 4,
drule-tac $x = N$ in spec,
drule-tac $x = N$ in spec, (erule exE)+,
rename-tac N M1 M2)
apply (subgoal-tac $\forall n > (\max M1 M2). (f n) \cdot_r (g n) \in (vp K v)(Vr K v)(an N),$
blast)
apply (rule allI, rule impI)
apply (drule-tac $x = M1$ and $y = M2$ in two-inequalities, assumption+,
thin-tac $\forall n > M2. g n \in vp K v (Vr K v)(an N)$)
apply (thin-tac $\forall j. f j \cdot_r g j \in \text{carrier } K,$
thin-tac $\forall j. f j \in \text{carrier } K,$
thin-tac $\forall j. g j \in \text{carrier } K,$
drule-tac $x = n$ in spec, simp, erule conjE,
erule conjE,
frule Vr-ring[of v])
apply (cut-tac $n = N$ in an-nat-pos)
apply (frule-tac $x = f n$ in mem-vp-apow-mem-Vr[of v], assumption+,
frule-tac $x = g n$ in mem-vp-apow-mem-Vr[of v], assumption+,
frule-tac $n = an N$ in vp-apow-ideal[of v], assumption+)
apply (frule-tac $I = vp K v (Vr K v)(an N)$ and $x = g n$ and
 $r = f n$ in Ring.ideal-ring-multiple[of Vr K v], assumption+,
simp add:Vr-tOp-f-tOp)

apply (rule allI)
apply (subgoal-tac $\forall j. (f j) \cdot_r (g j) =$
 $(f j) \cdot_r ((g j) \pm (-_a c)) \pm (f j) \cdot_r c,$ simp,
thin-tac $\forall j. (f j) \cdot_r (g j) =$
 $(f j) \cdot_r ((g j) \pm (-_a c)) \pm (f j) \cdot_r c,$
thin-tac $\forall j. (f j) \cdot_r (g j \pm -_a c) \pm (f j) \cdot_r c \in \text{carrier } K$)
apply (rotate-tac 4,
drule-tac $x = N + na (Abs (n-val K v c))$ in spec,
drule-tac $x = N$ in spec)
apply (erule exE)+ **apply** (rename-tac N M1 M2)
apply (subgoal-tac $\forall n. (\max M1 M2) < n \longrightarrow (f n) \cdot_r (g n \pm -_a c) \pm$
 $(f n) \cdot_r c \in vp K v (Vr K v)(an N),$ blast)

apply (*rule allI*, *rule impI*, *simp*, *erule conjE*,
drule-tac x = n in spec,
drule-tac x = n in spec,
drule-tac x = n in spec)
apply (*frule n-val-valuation*[of *v*])
apply (*frule value-in-aug-inf*[of *n-val K v c*], *assumption+*,
simp add:aug-inf-def)
apply (*frule val-nonzero-noninf*[of *n-val K v c*], *assumption+*)
apply (*cut-tac Abs-noninf*[of *n-val K v c*])
apply (*cut-tac Abs-pos*[of *n-val K v c*]) **apply** (*simp add:an-na*)

apply (*drule-tac x = n in spec*, *simp*)
apply (*frule-tac b = f n and n = an N in convergenceTr1*[of
v c], *assumption+*)
apply *simp*

apply (*frule-tac x = f n and N = an N + Abs (n-val K v c) in*
mem-vp-apow-mem-Vr[of *v*],
frule-tac n = an N in vp-apow-ideal[of *v*])
apply *simp*
apply (*rule-tac x = an N and y = Abs (n-val K v c) in aadd-two-pos*)
apply *simp* **apply** (*simp add:Abs-pos*) **apply** *assumption*

apply (*frule-tac x = g n $\pm$ ($-_a c$) and N = an N in mem-vp-apow-mem-Vr*[of
v], *simp*, *assumption+*) **apply** (
frule-tac x = c $\cdot_r$ (f n) and N = an N in mem-vp-apow-mem-Vr[of
v], *simp*) **apply** *simp*
apply (*simp add:Ring.ring-tOp-commute*[of *K c*]) **apply** (
frule Vr-ring[of *v*],
frule-tac I = (vp K v) (Vr K v) (an N) and x = g n $\pm$ ($-_a c$)
and r = f n in Ring.ideal-ring-multiple[of *Vr K v*])
apply (*simp add:vp-apow-ideal*) **apply** *assumption+*
apply (*frule-tac I = vp K v (Vr K v) (an N) and*
x = (f n) $\cdot_r$ (g n $\pm$ $-_a c$) and y = (f n) $\cdot_r$ c in
Ring.ideal-pOp-closed[of *Vr K v*])
apply (*simp add:vp-apow-ideal*)
apply (*simp add:Vr-tOp-f-tOp*, *assumption*)
apply (*simp add:Vr-tOp-f-tOp Vr-pOp-f-pOp*,
frule-tac x = (f n) $\cdot_r$ (g n $\pm$ $-_a c$) and N = an N in mem-vp-apow-mem-Vr[of
v], *simp add:Vr-pOp-f-pOp*, *assumption+*)
apply (*frule-tac N = an N and x = (f n) $\cdot_r$ c in mem-vp-apow-mem-Vr*[of
v]) **apply** *simp* **apply** *assumption*
apply (*simp add:Vr-pOp-f-pOp*) **apply** *simp*

apply (*thin-tac $\forall N. \exists M. \forall n > M. f n \in vp K v (Vr K v) (an N)$,*
thin-tac $\forall N. \exists M. \forall n > M. g n \pm -_a c \in vp K v (Vr K v) (an N)$,
rule allI)
apply (*drule-tac x = j in spec*,
drule-tac x = j in spec,

$\text{drule-tac } x = j \text{ in spec,}$
 $\text{frule aGroup.ag-mOp-closed[of K c], assumption+,}$
 $\text{frule-tac } x = f j \text{ in Ring.ring-tOp-closed[of K - c], assumption+,}$
 $\text{frule-tac } x = f j \text{ in Ring.ring-tOp-closed[of K - } -_a \text{ c], assumption+)}$
 $\text{apply (simp add:Ring.ring-distrib1, simp add:aGroup.ag-pOp-assoc,}$
 $\text{simp add:Ring.ring-distrib1 [THEN sym],}$
 $\text{simp add:aGroup.ag-l-inv1, simp add:Ring.ring-times-x-0,}$
 $\text{simp add:aGroup.ag-r-zero)}$
done

lemma (in Corps) $\text{limit-minus:} \llbracket \text{valuation } K v; \forall j. f j \in \text{carrier } K;$
 $b \in \text{carrier } K; \lim_{K v} f b \rrbracket \implies \lim_{K v} (\lambda j. (-_a (f j))) (-_a b)$
apply (simp add:limit-def)
apply (rule allI,
 $\text{rotate-tac -1, frule-tac } x = N \text{ in spec,}$
 $\text{thin-tac } \forall N. \exists M. \forall n. M < n \longrightarrow$
 $f n \pm -_a b \in (vp K v)(Vr K v) (an N),$
 erule exE,
 $\text{subgoal-tac } \forall n. M < n \longrightarrow$
 $(-_a (f n)) \pm (-_a (-_a b)) \in (vp K v)(Vr K v) (an N),$
 $\text{blast})$
apply (rule allI, rule impI,
 $\text{frule-tac } x = n \text{ in spec,}$
 $\text{frule-tac } x = n \text{ in spec, simp})$

apply (frule Vr-ring[of v],
 $\text{frule-tac } n = an N \text{ in vp-apow-ideal[of v], simp})$
apply (frule-tac $x = f n \pm -_a b$ and $N = an N$ in
 $\text{mem-vp-apow-mem-Vr[of v], simp+,}$
 $\text{frule-tac } I = vp K v (Vr K v) (an N) \text{ and } x = f n \pm (-_a b) \text{ in}$
 $\text{Ring.ideal-inv1-closed[of Vr K v], assumption+, simp})$
apply (simp add:Vr-mOp-f-mOp)
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
 $\text{frule aGroup.ag-mOp-closed[of K b], assumption+})$
apply (simp add:aGroup.ag-p-inv[of K])
done

lemma (in Corps) $\text{inv-diff:} \llbracket x \in \text{carrier } K; x \neq \mathbf{0}; y \in \text{carrier } K; y \neq \mathbf{0} \rrbracket \implies$
 $(x^{-K}) \pm (-_a (y^{-K})) = (x^{-K}) \cdot_r (y^{-K}) \cdot_r (-_a (x \pm (-_a y)))$
apply (cut-tac invf-closed1[of x], simp, erule conjE,
 $\text{cut-tac invf-closed1[of y], simp, erule conjE,}$
 $\text{cut-tac field-is-ring, frule Ring.ring-is-ag[of K],}$
 $\text{frule Ring.ring-tOp-closed[of K } x^{-K} y^{-K}], \text{ assumption+,}$
 $\text{frule aGroup.ag-mOp-closed[of K x], assumption+,}$
 $\text{frule aGroup.ag-mOp-closed[of K y], assumption+,}$
 $\text{frule aGroup.ag-mOp-closed[of K } x^{-K}], \text{ assumption+,}$
 $\text{frule aGroup.ag-mOp-closed[of K } y^{-K}], \text{ assumption+,}$
 $\text{frule aGroup.ag-pOp-closed[of K } x -_a y], \text{ assumption+})$

```

apply (simp add:Ring.ring-inv1-2[THEN sym],
      simp only:Ring.ring-distrib1[of K (x-K) ·r (y-K) x -a y],
      simp only:Ring.ring-tOp-commute[of K - x],
      simp only:Ring.ring-inv1-2[THEN sym, of K],
      simp only:Ring.ring-tOp-assoc[THEN sym],
      simp only:Ring.ring-tOp-commute[of K x],
      cut-tac linvf[of x], simp+,
      simp add:Ring.ring-l-one, simp only:Ring.ring-tOp-assoc,
      cut-tac linvf[of y], simp+,
      simp only:Ring.ring-r-one)
apply (simp add:aGroup.ag-p-inv[of K], simp add:aGroup.ag-inv-inv,
      rule aGroup.ag-pOp-commute[of K x-K -a y-K], assumption+)
apply simp+
done

lemma times2plus:(2::nat)*n = n + n
by simp

lemma (in Corps) limit-inv:⟦valuation K v; ∀j. f j ∈ carrier K;
  b ∈ carrier K; b ≠ 0; limK v f b⟧ ⇒
  limK v (λj. if (f j) = 0 then 0 else (f j)-K) (b-K)
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (frule limit-n-val[of b v], assumption+)
apply (subst limit-def)
apply (rule allI, erule exE)
apply (subgoal-tac ∀ m > Na. f m ≠ 0)
prefer 2
apply (rule allI, rule impI, rotate-tac -2,
  drule-tac x = m in spec, simp)
apply (frule n-val-valuation[of v])
apply (frule val-nonzero-noninf[of n-val K v b], assumption+)
apply (rule contrapos-pp, simp+, simp add:value-of-zero)
apply (unfold limit-def)
apply (rotate-tac 2,
  frule-tac x = N + 2*(na (Abs (n-val K v b))) in
  spec)
apply (erule exE)
apply (subgoal-tac ∀ n > (max Na M).
  (if f n = 0 then 0 else f n-K) ± -a b-K ∈ vp K v (Vr K v) (an N),
  blast)
apply (rule allI, rule impI)
apply (cut-tac x = Na and y = max Na M and z = n
  in le-less-trans)
apply simp+
apply (thin-tac ∀ N. ∃ M. ∀ n > M. f n ± -a b ∈ vp K v (Vr K v) (an N))
apply (drule-tac x = n in spec,
  drule-tac x = n in spec,
  drule-tac x = n in spec,
  drule-tac x = n in spec, simp)

```

apply (*subst inv-diff*, *assumption*+)
apply (*cut-tac* $x = f\ n$ **in** *invf-closed1*, *simp*,
 $cut-tac\ x = b$ **in** *invf-closed1*, *simp*, *simp*, (*erule conjE*)+)

apply (*frule-tac* $n = an\ N + an\ (2 * na\ (Abs\ (n-val\ K\ v\ b)))$ **and**
 $x = f\ n \pm -_a\ b$ **in** *n-value-x-1*[*of v*])
apply (*simp only:an-npn*[*THEN sym*], *rule an-nat-pos*)
apply *assumption*
apply (*rule-tac* $x = f\ n^- K \cdot_r b^- K \cdot_r (-_a (f\ n \pm -_a\ b))$ **and** $v = v$ **and**
 $n = an\ N$ **in** *n-value-x-2*, *assumption*+)

apply (*frule n-val-valuation*[*of v*])
apply (*subst val-pos-mem-Vr*[*THEN sym*, *of v*], *assumption*+)

apply (*rule Ring.ring-tOp-closed*, *assumption*+) +
apply (*rule aGroup.ag-mOp-closed*, *assumption*)
apply (*rule aGroup.ag-pOp-closed*, *assumption*+,
rule aGroup.ag-mOp-closed, *assumption*+)

apply (*subst val-pos-n-val-pos*[*of v*], *assumption*+,
rule Ring.ring-tOp-closed, *assumption*+,
rule Ring.ring-tOp-closed, *assumption*+,
rule aGroup.ag-mOp-closed, *assumption*+,
rule aGroup.ag-pOp-closed, *assumption*+,
rule aGroup.ag-mOp-closed, *assumption*+)

apply (*subst val-t2p*[*of n-val K v*], *assumption*+,
rule Ring.ring-tOp-closed, *assumption*+,
rule aGroup.ag-mOp-closed, *assumption*+,
rule aGroup.ag-pOp-closed, *assumption*+,
rule aGroup.ag-mOp-closed, *assumption*+,
subst val-minus-eq[*of n-val K v*], *assumption*+,
(*rule aGroup.ag-pOp-closed*, *assumption*+,
rule aGroup.ag-mOp-closed, *assumption*+)

apply (*subst val-t2p*[*of n-val K v*], *assumption*+)

apply (*simp add:value-of-inv*)

apply (*frule-tac* $x = an\ N + an\ (2 * na\ (Abs\ (n-val\ K\ v\ b)))$ **and** $y = n-val\ K$
 $v\ (f\ n \pm -_a\ b)$ **and** $z = -\ n-val\ K\ v\ b + -\ n-val\ K\ v\ b$ **in** *aadd-le-mono*)

apply (*cut-tac* $z = n-val\ K\ v\ b$ **in** *Abs-pos*)

apply (*frule val-nonzero-z*[*of n-val K v b*], *assumption*+, *erule exE*)

apply (*rotate-tac* -1 , *drule sym*, *cut-tac* $z = z$ **in** *z-neq-minf*,
 $cut-tac\ z = z$ **in** *z-neq-inf*, *simp*,
 $cut-tac\ a = (n-val\ K\ v\ b)$ **in** *Abs-noninf*, *simp*)

apply (*simp only:times2plus an-npn*, *simp add:an-na*)

apply (*rotate-tac* -4 , *drule sym*, *simp*)

apply (*thin-tac* $f\ n \pm -_a\ b \in vp\ K\ v\ (Vr\ K\ v)\ (an\ N + (ant\ |z| + ant\ |z|))$)

apply (*simp add:an-def*, *simp add:aminus*, (*simp add:a-zpz*)+)

apply (*subst aadd-commute*)

apply (*rule-tac* $i = 0$ **and** $j = ant\ (int\ N + 2 * |z| - 2 * z)$ **and**
 $k = n-val\ K\ v\ (f\ n \pm -_a\ b) + ant\ (-\ (2 * z))$ **in** *ale-trans*)

apply (*subst ant-0*[*THEN sym*])

apply (*subst ale-zle*, *simp*, *assumption*)

apply (*frule* *n-val-valuation*[*of v*])
apply (*subst val-t2p*[*of n-val K v*], *assumption*+)
apply (*rule* *Ring.ring-tOp-closed*, *assumption*+) +
apply (*rule* *aGroup.ag-mOp-closed*, *assumption*)
apply (*rule* *aGroup.ag-pOp-closed*, *assumption*+,
rule *aGroup.ag-mOp-closed*, *assumption*+)
apply (*subst val-t2p*[*of n-val K v*], *assumption*+)
apply (*subst val-minus-eq*[*of n-val K v*], *assumption*+,
rule *aGroup.ag-pOp-closed*, *assumption*+,
rule *aGroup.ag-mOp-closed*, *assumption*+)

apply (*simp add:value-of-inv*)
apply (*frule-tac* *x = an N + an (2 * na (Abs (n-val K v b)))* **and** *y = n-val K*
v (f n ± -_a b) **and** *z = - n-val K v b + - n-val K v b* **in** *aadd-le-mono*)
apply (*cut-tac* *z = n-val K v b* **in** *Abs-pos*)
apply (*frule val-nonzero-z*[*of n-val K v b*], *assumption*+, *erule exE*)
apply (*rotate-tac* -1, *drule sym*, *cut-tac* *z = z* **in** *z-neq-minf*,
cut-tac *z = z* **in** *z-neq-inf*, *simp*,
cut-tac *a = (n-val K v b)* **in** *Abs-noninf*, *simp*)
apply (*simp only:times2plus an-npn*, *simp add:an-na*)
apply (*rotate-tac* -4, *drule sym*, *simp*)
apply (*thin-tac* *f n ± -_a b ∈ vp K v (Vr K v) (an N + (ant |z| + ant |z|))*)
apply (*simp add:an-def*, *simp add:aminus*, (*simp add:a-zpz*) +)
apply (*subst aadd-commute*)
apply (*rule-tac* *i = ant (int N)* **and** *j = ant (int N + 2 * |z| - 2 * z)*
and *k = n-val K v (f n ± -_a b) + ant (- (2 * z))* **in** *ale-trans*)
apply (*subst ale-zle*, *simp*, *assumption*)

apply *simp*
done

definition

Cauchy-seq :: [*-*, '*b* ⇒ *ant*, *nat* ⇒ '*b*]
⇒ *bool* (⟨(3*Cauchy* - -)⟩ [90,90,91]90) **where**
Cauchy_K v f ⇔ (∀ *n*. (*f n*) ∈ *carrier K*) ∧ (
∀ *N*. ∃ *M*. (∀ *n m*. *M* < *n* ∧ *M* < *m* →
((*f n*) ±_{*K*} (-_{*a*}*K* (*f m*))) ∈ (*vp K v*)(*Vr K v*) (*an N*)))

definition

v-complete :: [*'b* ⇒ *ant*, -] ⇒ *bool*
(⟨(2*Complete* - -)⟩ [90,91]90) **where**
Complete_v K ⇔ (∀ *f*. (*Cauchy_K v f*) →
(∃ *b*. *b* ∈ (*carrier K*) ∧ *lim_K v f b*))

lemma (**in** *Corps*) *has-limit-Cauchy*: [valuation *K v*; ∀ *j*. *f j* ∈ *carrier K*;
b ∈ *carrier K*; *lim_K v f b*] ⇒ *Cauchy_K v f*
apply (*simp add:Cauchy-seq-def*)
apply (*rule allI*)
apply (*simp add:limit-def*)

```

apply (rotate-tac -1)
apply (drule-tac x = N in spec)
apply (erule exE)
apply (subgoal-tac  $\forall n\ m. M < n \wedge M < m \longrightarrow$ 
 $f\ n \pm -_a (f\ m) \in vp\ K\ v^{(Vr\ K\ v)} (an\ N)$ )

apply blast
apply ((rule allI)+, rule impI, erule conjE)
apply (frule-tac x = n in spec,
  frule-tac x = m in spec,
  thin-tac  $\forall j. f\ j \in carrier\ K$ ,
  frule-tac x = n in spec,
  frule-tac x = m in spec,
  thin-tac  $\forall n. M < n \longrightarrow f\ n \pm -_a b \in vp\ K\ v^{(Vr\ K\ v)} (an\ N)$ ,
  simp)
apply (frule-tac n = an N in vp-apow-ideal[of v], simp)
apply (frule Vr-ring[of v])
apply (frule-tac x = f m  $\pm -_a b$  and I = vp K v(Vr K v) (an N) in
  Ring.ideal-inv1-closed[of Vr K v], assumption+)
apply (frule-tac h = f m  $\pm -_a b$  and I = vp K v(Vr K v) (an N) in
  Ring.ideal-subset[of Vr K v], assumption+,
  frule-tac h = f n  $\pm -_a b$  and I = vp K v(Vr K v) (an N) in
  Ring.ideal-subset[of Vr K v], assumption+)
apply (frule-tac h = -a Vr K v (f m  $\pm -_a b$ ) and I = vp K v(Vr K v) (an N) in
  Ring.ideal-subset[of Vr K v], assumption+,
  frule-tac h = f n  $\pm -_a b$  and I = vp K v(Vr K v) (an N) in
  Ring.ideal-subset[of Vr K v], assumption+)
apply (frule-tac I = (vp K v) (Vr K v) (an N) and x = f n  $\pm -_a b$  and
  y = -a (Vr K v) (f m  $\pm -_a b$ ) in Ring.ideal-pOp-closed[of Vr K v],
  assumption+)
apply (simp add: Vr-pOp-f-pOp) apply (simp add: Vr-mOp-f-mOp)
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (frule aGroup.ag-mOp-closed[of K b], assumption+)
apply (frule-tac x = f m  $\pm -_a b$  in Vr-mem-f-mem[of v], assumption+)
apply (frule-tac x = f m  $\pm -_a b$  in aGroup.ag-mOp-closed[of K],
  assumption+)
apply (simp add: aGroup.ag-pOp-assoc)
apply (simp add: aGroup.ag-pOp-commute[of K -a b])
apply (simp add: aGroup.ag-p-inv[of K])
apply (frule-tac x = f m in aGroup.ag-mOp-closed[of K], assumption+)
apply (simp add: aGroup.ag-inv-inv)
apply (simp add: aGroup.ag-pOp-assoc[of K - b -a b])
apply (simp add: aGroup.ag-r-inv1[of K], simp add: aGroup.ag-r-zero)
done

lemma (in Corps) no-limit-zero-Cauchy:  $\llbracket valuation\ K\ v; Cauchy_{K\ v}\ f;$ 
 $\neg (\lim_{K\ v}\ f\ \mathbf{0}) \rrbracket \Longrightarrow$ 
 $\exists N\ M. (\forall m. N < m \longrightarrow ((n-val\ K\ v)\ (f\ M)) = ((n-val\ K\ v)\ (f\ m)))$ 
apply (frule not-limit-zero[of v f], thin-tac  $\neg \lim_{K\ v}\ f\ \mathbf{0}$ )

```

apply (*simp add:Cauchy-seq-def*, *assumption*) **apply** (*erule exE*)
apply (*rename-tac L*)
apply (*simp add:Cauchy-seq-def*, *erule conjE*,
rotate-tac -1,
frule-tac x = L in spec, *thin-tac $\forall N. \exists M. \forall n m.$*
 $M < n \wedge M < m \longrightarrow f n \pm -_a (f m) \in vp K v^{(Vr K v) (an N)}$)
apply (*erule exE*)
apply (*drule-tac x = M in spec*)
apply (*erule exE*, *erule conjE*)
apply (*rotate-tac -3*,
frule-tac x = m in spec)
apply (*thin-tac $\forall n m. M < n \wedge M < m \longrightarrow$*
 $f n \pm -_a (f m) \in (vp K v)^{(Vr K v) (an L)}$)
apply (*subgoal-tac $M < m \wedge (\forall ma. M < ma \longrightarrow$*
 $n-val K v (f m) = n-val K v (f ma))$)
apply *blast*
apply *simp*

apply (*rule allI*, *rule impI*)
apply (*rotate-tac -2*)
apply (*drule-tac x = ma in spec*)
apply *simp*

apply (*frule Vr-ring[of v]*,
frule-tac n = an L in vp-apow-ideal[of v], *simp*)
apply (*frule-tac $I = vp K v^{(Vr K v) (an L)}$* **and** *$x = f m \pm -_a (f ma)$*
in Ring.ideal-inv1-closed[of Vr K v], *assumption+*) **apply** (
frule-tac $I = vp K v^{(Vr K v) (an L)}$ **and**
 $h = f m \pm -_a (f ma)$ in Ring.ideal-subset[of Vr K v],
assumption+, *simp add:Vr-mOp-f-mOp*)
apply (*frule-tac x = m in spec*,
drule-tac x = ma in spec) **apply** (
cut-tac field-is-ring, *frule Ring.ring-is-ag[of K]*,
frule-tac x = f ma in aGroup.ag-mOp-closed[of K], *assumption+*,
frule-tac x = f m and y = -_a (f ma) in aGroup.ag-p-inv[of K],
assumption+, *simp only:aGroup.ag-inv-inv*,
frule-tac x = f m in aGroup.ag-mOp-closed[of K], *assumption+*,
simp add:aGroup.ag-pOp-commute,
thin-tac -_a (f m $\pm$ -_a (f ma)) = f ma $\pm$ -_a (f m),
thin-tac f m $\pm$ -_a (f ma) $\in$ vp K v^{(Vr K v) (an L)})

apply (*frule-tac x = f ma $\pm$ -_a (f m) and n = an L in n-value-x-1[of*
v], *simp*) **apply** *assumption* **apply** (
frule n-val-valuation[of v],
frule-tac x = f m and y = f ma $\pm$ -_a (f m) in value-less-eq[of
n-val K v], *assumption+*) **apply** (*simp add:aGroup.ag-pOp-closed*)
apply (

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    rule-tac x = n-val K v (f m) and y = an L and
    z = n-val K v (f ma ± -a (f m)) in
    aless-le-trans, assumption+)
  apply (frule-tac x = f ma ± -a (f m) in Vr-mem-f-mem[of v])
  apply (simp add:Ring.ideal-subset)
  apply (frule-tac x = f m and y = f ma ± -a (f m) in
    aGroup.ag-pOp-commute[of K], assumption+)
  apply (simp add:aGroup.ag-pOp-assoc,
    simp add:aGroup.ag-l-inv1, simp add:aGroup.ag-r-zero)
done

lemma (in Corps) no-limit-zero-Cauchy1: [valuation K v; ∀ j. f j ∈ carrier K;
  CauchyK v f; ¬ (limK v f 0)] ⇒ ∃ N M. (∀ m. N < m → v (f M) = v (f m))
  apply (frule no-limit-zero-Cauchy[of v f], assumption+)
  apply (erule exE)+
  apply (subgoal-tac ∀ m. N < m → v (f M) = v (f m)) apply blast
  apply (rule allI, rule impI)
  apply (frule-tac x = M in spec,
    drule-tac x = m in spec,
    drule-tac x = m in spec, simp)
  apply (simp add:n-val[THEN sym, of v])
done

definition
  subfield :: [-, ('b, 'm1) Ring-scheme] ⇒ bool where
  subfield K K' ⟷ Corps K' ∧ carrier K ⊆ carrier K' ∧
    idmap (carrier K) ∈ rHom K K'

definition
  v-completion :: ['b ⇒ ant, 'b ⇒ ant, -, ('b, 'm) Ring-scheme] ⇒ bool
  (⟨(4 Completion - - -)⟩ [90,90,90,91] 90) where
  Completionv v' K K' ⟷ subfield K K' ∧
    Completev v' K' ∧ (∀ x ∈ carrier K. v x = v' x) ∧
    (∀ x ∈ carrier K'. (∃ f. CauchyK v f ∧ limK' v' f x))

lemma (in Corps) subfield-zero: [Corps K'; subfield K K'] ⇒ 0K = 0K'
  apply (simp add:subfield-def, (erule conjE)+)
  apply (simp add:rHom-def, (erule conjE)+)
  apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
    frule Corps.field-is-ring[of K'], frule Ring.ring-is-ag[of K'])
  apply (frule aHom-0-0[of K K' IK], assumption+)
  apply (frule aGroup.ag-inc-zero[of K], simp add:idmap-def)
done

lemma (in Corps) subfield-pOp: [Corps K'; subfield K K'; x ∈ carrier K;
  y ∈ carrier K] ⇒ x ± y = x ±K' y
  apply (cut-tac field-is-ring, frule Corps.field-is-ring[of K'],
    frule Ring.ring-is-ag[of K], frule Ring.ring-is-ag[of K'])
  apply (simp add:subfield-def, erule conjE, simp add:rHom-def,

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      frule conjunct1)
apply (thin-tac  $I_K \in aHom\ K\ K' \wedge$ 
       $(\forall x \in carrier\ K. \forall y \in carrier\ K. I_K\ (x \cdot_r y) = I_K\ x \cdot_r_{K'} I_K\ y) \wedge$ 
       $I_K\ 1_r = 1_{r\ K'})$ 
apply (frule aHom-add[of  $K\ K'\ I_K\ x\ y$ ], assumption+,
      frule aGroup.ag-pOp-closed[of  $K\ x\ y$ ], assumption+,
      simp add:idmap-def)
done

lemma (in Corps) subfield-mOp:  $\llbracket Corps\ K';\ subfield\ K\ K';\ x \in carrier\ K \rrbracket \implies$ 
       $\neg_a\ x = \neg_{a\ K'}\ x$ 
apply (cut-tac field-is-ring, frule Corps.field-is-ring[of  $K^\intercal$ ],
      frule Ring.ring-is-ag[of  $K$ ], frule Ring.ring-is-ag[of  $K^\intercal$ ])
apply (simp add:subfield-def, erule conjE, simp add:rHom-def,
      frule conjunct1)
apply (thin-tac  $I_K \in aHom\ K\ K' \wedge$ 
       $(\forall x \in carrier\ K. \forall y \in carrier\ K. I_K\ (x \cdot_r y) = I_K\ x \cdot_r_{K'} I_K\ y) \wedge$ 
       $I_K\ 1_r = 1_{r\ K'})$ 
apply (frule aHom-inv-inv[of  $K\ K'\ I_K\ x$ ], assumption+,
      frule aGroup.ag-mOp-closed[of  $K\ x$ ], assumption+)
apply (simp add:idmap-def)
done

lemma (in Corps) completion-val-eq:  $\llbracket Corps\ K';\ valuation\ K\ v;\ valuation\ K'\ v';$ 
       $x \in carrier\ K;\ Completion_{v\ v'}\ K\ K^\intercal \rrbracket \implies v\ x = v'\ x$ 
apply (unfold v-completion-def, (erule conjE)+)
apply simp
done

lemma (in Corps) completion-subset:  $\llbracket Corps\ K';\ valuation\ K\ v;\ valuation\ K'\ v';$ 
       $Completion_{v\ v'}\ K\ K^\intercal \rrbracket \implies carrier\ K \subseteq carrier\ K'$ 
apply (unfold v-completion-def, (erule conjE)+)
apply (simp add:subfield-def)
done

lemma (in Corps) completion-subfield:  $\llbracket Corps\ K';\ valuation\ K\ v;$ 
       $valuation\ K'\ v';\ Completion_{v\ v'}\ K\ K^\intercal \rrbracket \implies subfield\ K\ K'$ 
apply (simp add:v-completion-def)
done

lemma (in Corps) subfield-sub:  $subfield\ K\ K' \implies carrier\ K \subseteq carrier\ K'$ 
apply (simp add:subfield-def)
done

lemma (in Corps) completion-Vring-sub:  $\llbracket Corps\ K';\ valuation\ K\ v;$ 
       $valuation\ K'\ v';\ Completion_{v\ v'}\ K\ K^\intercal \rrbracket \implies$ 
       $carrier\ (Vr\ K\ v) \subseteq carrier\ (Vr\ K'\ v')$ 
apply (rule subsetI,
      frule completion-subset[of  $K'\ v\ v^\intercal$ ], assumption+,

```

```

    frule-tac x = x in Vr-mem-f-mem[of v], assumption+,
    frule-tac x = x in completion-val-eq[of K' v v'],
    assumption+)
apply (frule-tac x1 = x in val-pos-mem-Vr[THEN sym, of v],
    assumption+, simp,
    frule-tac c = x in subsetD[of carrier K carrier K'], assumption+,
    simp add:Corps.val-pos-mem-Vr[of K' v'])
done

lemma (in Corps) completion-idmap-rHom:[Corps K'; valuation K v;
    valuation K' v'; Completionv v' K K']  $\implies$ 
    I(Vr K v)  $\in$  rHom (Vr K v) (Vr K' v')
apply (frule completion-Vring-sub[of K' v v'],
    assumption+,
    frule completion-subfield[of K' v v'],
    assumption+,
    frule Vr-ring[of v],
    frule Ring.ring-is-ag[of Vr K v],
    frule Corps.Vr-ring[of K' v'], assumption+,
    frule Ring.ring-is-ag[of Vr K' v'])
apply (simp add:rHom-def)
apply (rule conjI)
apply (simp add:aHom-def,
    rule conjI,
    simp add:idmap-def, simp add:subsetD)
apply (rule conjI)
apply (simp add:idmap-def extensional-def)
apply ((rule ballI)+) apply (
    frule-tac x = a and y = b in aGroup.ag-pOp-closed, assumption+,
    simp add:idmap-def,
    frule-tac c = a in subsetD[of carrier (Vr K v)
        carrier (Vr K' v')], assumption+,
    frule-tac c = b in subsetD[of carrier (Vr K v)
        carrier (Vr K' v')], assumption+,
    simp add:Vr-pOp-f-pOp,
    frule-tac x = a in Vr-mem-f-mem[of v], assumption,
    frule-tac x = b in Vr-mem-f-mem[of v], assumption,
    simp add:Corps.Vr-pOp-f-pOp,
    simp add:subfield-pOp)
apply (rule conjI)
apply ((rule ballI)+,
    frule-tac x = x and y = y in Ring.ring-tOp-closed, assumption+,
    simp add:idmap-def, simp add:subfield-def, erule conjE)
apply (frule-tac c = x in subsetD[of carrier (Vr K v)
    carrier (Vr K' v')], assumption+,
    frule-tac c = y in subsetD[of carrier (Vr K v)
    carrier (Vr K' v')], assumption+)
apply (simp add:Vr-tOp-f-tOp Corps.Vr-tOp-f-tOp)
apply (frule-tac x = x in Vr-mem-f-mem[of v], assumption+,

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    frule-tac x = y in Vr-mem-f-mem[of v], assumption+,
    frule-tac x = x in Corps.Vr-mem-f-mem[of K' v], assumption+,
    frule-tac x = y in Corps.Vr-mem-f-mem[of K' v], assumption+,
    cut-tac field-is-ring, frule Corps.field-is-ring[of K],
    frule-tac x = x and y = y in Ring.ring-top-closed[of K], assumption+)
  apply (frule-tac x = x and y = y in rHom-top[of K K' - I_K],
    assumption+, simp add:idmap-def)
  apply (frule Ring.ring-one[of Vr K v], simp add:idmap-def)
  apply (simp add:Vr-1-f-1 Corps.Vr-1-f-1)
  apply (simp add:subfield-def, (erule conjE)+)
  apply (cut-tac field-is-ring, frule Corps.field-is-ring[of K],
    frule Ring.ring-one[of K],
    frule rHom-one[of K K' I_K], assumption+, simp add:idmap-def)
done

```

```

lemma (in Corps) completion-vpr-sub: [Corps K'; valuation K v; valuation K' v';
  Completionv v' K K']  $\implies$  vp K v  $\subseteq$  vp K' v'
apply (rule subsetI,
  frule completion-subset[of K' v v'], assumption+,
  frule Vr-ring[of v], frule vp-ideal[of v],
  frule-tac h = x in Ring.ideal-subset[of Vr K v vp K v],
  assumption+,
  frule-tac x = x in Vr-mem-f-mem[of v], assumption+,
  frule-tac x = x in completion-val-eq[of K' v v'],
  assumption+)
apply (frule completion-subset[of K' v v'],
  assumption+,
  frule-tac c = x in subsetD[of carrier K carrier K'], assumption+,
  simp add:Corps.vp-mem-val-poss vp-mem-val-poss)
done

```

```

lemma (in Corps) val-v-completion: [Corps K'; valuation K v; valuation K' v';
  x  $\in$  carrier K'; x  $\neq$  0K'; Completionv v' K K']  $\implies$ 
   $\exists f. (Cauchy_K v f) \wedge (\exists N. (\forall m. N < m \longrightarrow v (f m) = v' x))$ 
apply (simp add:v-completion-def, erule conjE, (erule conjE)+)
apply (rotate-tac -1, drule-tac x = x in bspec, assumption+,
  erule exE, erule conjE,
  subgoal-tac  $\exists N. \forall m. N < m \longrightarrow v (f m) = v' x$ , blast)
  thm Corps.limit-val
apply (frule-tac f = f and v = v' in Corps.limit-val[of K' x],
  assumption+,
  unfold Cauchy-seq-def, frule conjunct1, fold Cauchy-seq-def)
apply (rule allI, drule-tac x = j in spec,
  simp add:subfield-def, erule conjE, simp add:subsetD, assumption+)
apply (simp add:Cauchy-seq-def)
done

```

```

lemma (in Corps) v-completion-v-limit: [Corps K'; valuation K v;
  x  $\in$  carrier K; subfield K K'; Completev K';  $\forall j. f j \in$  carrier K;

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$\text{valuation } K' v'; \forall x \in \text{carrier } K. v x = v' x; \lim_{K' v'} f x \implies \lim_{K v} f x$
apply (*cut-tac field-is-ring*, *frule Ring.ring-is-ag*[*of K*],
frule Corps.field-is-ring[*of K'*], *frule Ring.ring-is-ag*[*of K'*],
subgoal-tac $\forall j. f j \in \text{carrier } K'$,
unfold subfield-def, *frule conjunct2*, *fold subfield-def*, *erule conjE*)
apply (*frule subsetD*[*of carrier K carrier K' x*], *assumption+*,
simp add:limit-diff-val[*of x f*],
subgoal-tac $\forall n. f n \pm_{K'} -_a K' x = f n \pm -_a x$)
apply (*rule allI*)
apply (*simp add:limit-def*)
apply (*rotate-tac 6*, *drule-tac* $x = N$ **in** *spec*,
erule exE)
apply (*subgoal-tac* $\forall n > M. \text{an } N \leq v'(f n \pm -_a x)$,
subgoal-tac $\forall n. (v'(f n \pm -_a x) = v(f n \pm -_a x))$, *simp*, *blast*)
apply (*rule allI*)
apply (*frule-tac* $x = f n \pm -_a x$ **in** *bspec*,
rule aGroup.ag-pOp-closed, *assumption+*, *simp*,
rule aGroup.ag-mOp-closed, *assumption+*) **apply** *simp*
apply (*rule allI*, *rule impI*)
apply (*frule-tac* $v = v'$ **and** $n = \text{an } N$ **and** $x = f n \pm -_a x$ **in**
Corps.n-value-x-1[*of K'*], *assumption+*, *simp*, *simp*)
apply (*frule-tac* $v = v'$ **and** $x = f n \pm -_a x$ **in** *Corps.n-val-le-val*[*of K'*],
assumption+)
apply (*cut-tac* $x = f n$ **and** $y = -_a x$ **in** *aGroup.ag-pOp-closed*, *assumption*,
simp, *rule aGroup.ag-mOp-closed*, *assumption+*, *simp add:subsetD*)
apply (*subst Corps.val-pos-n-val-pos*[*of K' v'*], *assumption+*)
apply (*cut-tac* $x = f n$ **and** $y = -_a x$ **in** *aGroup.ag-pOp-closed*, *assumption*,
simp, *rule aGroup.ag-mOp-closed*, *assumption+*, *simp add:subsetD*)
apply (*rule-tac* $i = 0$ **and** $j = \text{an } N$ **and** $k = \text{n-val } K' v' (f n \pm -_a x)$ **in**
ale-trans, *simp+*, *rule allI*)
apply (*subst subfield-pOp*[*of K'*], *assumption+*, *simp+*,
rule aGroup.ag-mOp-closed, *assumption+*)
apply (*simp add:subfield-mOp*[*of K'*])
apply (*cut-tac subfield-sub*[*of K'*], *simp add:subsetD*, *assumption+*)
done

lemma (**in** *Corps*) *Vr-idmap-aHom*: $\llbracket \text{Corps } K'; \text{valuation } K v; \text{valuation } K' v';$
 $\text{subfield } K K'; \forall x \in \text{carrier } K. v x = v' x \rrbracket \implies$
 $I_{(Vr K v)} \in \text{aHom } (Vr K v) (Vr K' v')$
apply (*simp add:aHom-def*)
apply (*subgoal-tac* $I_{(Vr K v)} \in \text{carrier } (Vr K v) \rightarrow \text{carrier } (Vr K' v')$)
apply *simp*
apply (*rule conjI*)
apply (*simp add:idmap-def*)
apply (*rule ballI*)
apply (*frule Vr-ring*[*of v*],
frule Ring.ring-is-ag[*of Vr K v*],
frule Corps.Vr-ring[*of K' v'*], *assumption+*,
frule Ring.ring-is-ag[*of Vr K' v'*])

apply (*frule-tac* $x = a$ **and** $y = b$ **in** *aGroup.ag-pOp-closed*[*of* $Vr\ K\ v$],
assumption+,
frule-tac $x = a$ **in** *funcset-mem*[*of* $I(Vr\ K\ v)$
carrier ($Vr\ K\ v$) *carrier* ($Vr\ K'\ v'$)], *assumption+*,
frule-tac $x = b$ **in** *funcset-mem*[*of* $I(Vr\ K\ v)$
carrier ($Vr\ K\ v$) *carrier* ($Vr\ K'\ v'$)], *assumption+*,
frule-tac $x = (I(Vr\ K\ v))\ a$ **and** $y = (I(Vr\ K\ v))\ b$ **in**
aGroup.ag-pOp-closed[*of* $Vr\ K'\ v'$], *assumption+*,
simp add: Vr-pOp-f-pOp)
apply (*simp add: idmap-def*, *simp add: subfield-def*, *erule conjE*,
simp add: rHom-def, *frule conjunct1*,
thin-tac $I_K \in aHom\ K\ K' \wedge$
 $(\forall x \in carrier\ K. \forall y \in carrier\ K. I_K\ (x \cdot_r y) = I_K\ x \cdot_{rK'} I_K\ y) \wedge$
 $I_K\ 1_r = 1_{rK'})$
apply (*simp add: Corps.Vr-pOp-f-pOp*[*of* $K'\ v'$])
apply (*frule-tac* $x = a$ **in** *Vr-mem-f-mem*[*of* v], *assumption+*,
frule-tac $x = b$ **in** *Vr-mem-f-mem*[*of* v], *assumption+*)
apply (*cut-tac field-is-ring*, *frule Ring.ring-is-ag*[*of* K],
frule Corps.field-is-ring[*of* K'], *frule Ring.ring-is-ag*[*of* K'])
apply (*frule-tac* $a = a$ **and** $b = b$ **in** *aHom-add*[*of* $K\ K'\ I_K$], *assumption+*,
frule-tac $x = a$ **and** $y = b$ **in** *aGroup.ag-pOp-closed*[*of* K], *assumption+*,
simp add: idmap-def)
apply (*rule Pi-I*,
drule-tac $x = x$ **in** *bspec*, *simp add: Vr-mem-f-mem*)
apply (*simp add: idmap-def*)
apply (*frule-tac* $x1 = x$ **in** *val-pos-mem-Vr*[*THEN sym*, *of* v],
simp add: Vr-mem-f-mem, *simp*)
apply (*subst Corps.val-pos-mem-Vr*[*THEN sym*, *of* $K'\ v'$], *assumption+*,
frule-tac $x = x$ **in** *Vr-mem-f-mem*[*of* v], *assumption+*,
frule subfield-sub[*of* K'], *simp add: subsetD*)
apply *assumption*
done

lemma *amult-pos-pos*: $0 \leq a \implies 0 \leq a * an\ N$
apply (*case-tac* $N = 0$, *simp add: an-0*)
apply (*case-tac* $a = \infty$, *simp*)
apply (*frule apos-neq-minf*[*of* a])
apply (*subst ant-tna*[*THEN sym*, *of* a], *simp*)
apply (*subst amult-0-r*, *simp*)
apply (*case-tac* $a = \infty$, *simp add: an-def*)
apply (*frule apos-neq-minf*[*of* a])
apply (*subst ant-tna*[*THEN sym*, *of* a], *simp*)
apply (*case-tac* $a = 0$, *simp*)
apply (*simp add: ant-0 an-def amult-0-l*)
apply (*cut-tac amult-pos1*[*of* $tna\ a\ an\ N$])
apply (*simp add: ant-tna*)
apply (*rule-tac ale-trans*[*of* $0\ an\ N\ a * an\ N$], *simp+*)
apply (*frule ale-neq-less*[*of* $0\ a$], *rule not-sym*, *assumption*)
apply (*subst aless-zless*[*THEN sym*, *of* $0\ tna\ a$], *simp add: ant-tna ant-0*)

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apply simp
done

lemma (in Corps) Cauchy-down: $\llbracket$ Corps  $K'$ ; valuation  $K$   $v$ ; valuation  $K'$   $v'$ ;
  subfield  $K$   $K'$ ;  $\forall x \in \text{carrier } K. v\ x = v'\ x$ ;  $\forall j. f\ j \in \text{carrier } K$ ; Cauchy $_{K' v'}$   $f$  $\rrbracket$ 
 $\implies$  Cauchy $_{K v}$   $f$ 
apply (simp add:Cauchy-seq-def, rule allI, erule conjE)
apply (rotate-tac  $-1$ , drule-tac
   $x = na\ (Lv\ K\ v) * N$  in spec,
  erule exE,
  subgoal-tac  $\forall n\ m. M < n \wedge M < m \longrightarrow$ 
     $f\ n \pm (-_a\ (f\ m)) \in vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$ , blast)
apply (simp add:amult-an-an) apply (simp add:an-na-Lv)
apply ((rule allI)+, rule impI, erule conjE) apply (
  rotate-tac  $-3$ , drule-tac  $x = n$  in spec,
  rotate-tac  $-1$ , drule-tac  $x = m$  in spec,
  simp)
apply (rotate-tac  $7$ ,
  frule-tac  $x = n$  in spec,
  drule-tac  $x = m$  in spec)
apply (simp add:subfield-mOp[THEN sym],
  cut-tac field-is-ring, frule Ring.ring-is-ag[of  $K$ ],
  frule-tac  $x = f\ m$  in aGroup.ag-mOp-closed[of  $K$ ], assumption+)
apply (simp add:subfield-pOp[THEN sym])
apply (frule-tac  $x = f\ n$  and  $y = -_a\ f\ m$  in aGroup.ag-pOp-closed,
  assumption+,
  frule subfield-sub[of  $K^\eta$ ],
  frule-tac  $c = f\ n \pm -_a\ f\ m$  in subsetD[of carrier  $K$  carrier  $K^\eta$ ],
  assumption+)
apply (frule Lv-pos[of  $v$ ],
  frule aless-imp-le[of  $0\ Lv\ K\ v$ ])
apply (frule-tac  $N = N$  in amult-pos-pos[of  $Lv\ K\ v$ ])
apply (frule-tac  $n = (Lv\ K\ v) * an\ N$  and  $x = f\ n \pm -_a\ f\ m$  in
  Corps.n-value-x-1[of  $K'\ v'$ ], assumption+)
apply (frule-tac  $x = f\ n \pm -_a\ f\ m$  in Corps.n-val-le-val[of  $K'\ v'$ ],
  assumption+,
  frule-tac  $j = Lv\ K\ v * an\ N$  and  $k = n\text{-val}\ K'\ v'\ (f\ n \pm -_a\ f\ m)$  in
  ale-trans[of  $0$ ], assumption+, simp add:Corps.val-pos-n-val-pos)
apply (frule-tac  $i = Lv\ K\ v * an\ N$  and  $j = n\text{-val}\ K'\ v'\ (f\ n \pm -_a\ f\ m)$ 
  and  $k = v'\ (f\ n \pm -_a\ f\ m)$  in ale-trans, assumption+,
  thin-tac  $n\text{-val}\ K'\ v'\ (f\ n \pm -_a\ f\ m) \leq v'\ (f\ n \pm -_a\ f\ m)$ ,
  thin-tac  $Lv\ K\ v * an\ N \leq n\text{-val}\ K'\ v'\ (f\ n \pm -_a\ f\ m)$ )
apply (rotate-tac  $1$ ,
  drule-tac  $x = f\ n \pm -_a\ f\ m$  in bspec, assumption,
  rotate-tac  $-1$ , drule sym, simp)
apply (frule-tac  $v1 = v$  and  $x1 = f\ n \pm -_a\ f\ m$  in n-val[THEN sym],
  assumption)
apply simp
apply (simp only:amult-commute[of  $-Lv\ K\ v$ ])

```

```

apply (frule Lv-z[of v], erule exE)

apply (cut-tac w = z and x = an N and y = n-val K v (f n  $\pm$   $-_a$  f m) in
  amult-pos-mono-l,
  cut-tac m = 0 and n = z in aless-zless, simp add:ant-0)
apply simp
apply (rule-tac x=f n  $\pm$   $-_a$  (f m) and n = an N in n-value-x-2[of v],
  assumption+)
apply (subst val-pos-mem-Vr[THEN sym, of v], assumption+)
apply (subst val-pos-n-val-pos[of v], assumption+)
apply (rule-tac j = an N and k = n-val K v (f n  $\pm$   $-_a$  f m) in
  ale-trans[of 0], simp, assumption+)
apply simp
done

lemma (in Corps) Cauchy-up:[Corps K'; valuation K v; valuation K' v';
  Completionv v' K K'; CauchyK v f]  $\implies$  CauchyK' v' f
apply (simp add:Cauchy-seq-def,
  erule conjE,
  rule conjI, unfold v-completion-def, frule conjunct1,
  fold v-completion-def, rule allI, frule subfield-sub[of K'])
apply (simp add:subsetD)

apply (rule allI)
apply (rotate-tac -1, drule-tac x = na (Lv K' v') * N
  in spec, erule exE)
apply (subgoal-tac  $\forall n m. M < n \wedge M < m \longrightarrow$ 
  f n  $\pm_{K'} (-_a_{K'} (f m)) \in vp K' v' (Vr K' v') (an N)$ , blast,
  (rule allI)+, rule impI, erule conjE)
apply (rotate-tac -3, drule-tac x = n in spec,
  rotate-tac -1,
  drule-tac x = m in spec, simp,
  frule-tac x = n in spec,
  drule-tac x = m in spec)
apply(unfold v-completion-def, frule conjunct1, fold v-completion-def,
  cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
  frule-tac x = f m in aGroup.ag-mOp-closed[of K], assumption+,
  frule-tac x = f n and y =  $-_a$  (f m) in aGroup.ag-pOp-closed[of K],
  assumption+)
apply (simp add:amult-an-an) apply (simp add:Corps.an-na-Lv)
apply (simp add:subfield-mOp, simp add:subfield-pOp) apply (
  frule-tac x = f n  $\pm_{K'} -_a_{K'} f m$  and n = (Lv K' v') * (an N)
  in n-value-x-1[of v])
apply (frule Corps.Lv-pos[of K' v'], assumption+,
  frule Corps.Lv-z[of K' v'],
  assumption, erule exE, simp,
  cut-tac n = N in an-nat-pos,
  frule-tac w1 = z and x1 = 0 and y1 = an N in
  amult-pos-mono-l[THEN sym], simp, simp add:amult-0-r)

```

apply *assumption*
apply (*frule-tac* $x = f\ n \pm_{K'} -_{a_{K'}} f\ m$ **in** *n-val-le-val*[*of v*],
assumption+)

apply (*subst* *n-val*[*THEN sym*, *of v*], *assumption*+)

apply (*frule* *Lv-pos*[*of v*], *frule* *Lv-z*[*of v*], *erule* *exE*, *simp*)

apply (*frule* *Corps.Lv-pos*[*of K' v*], *assumption*+,
frule *Corps.Lv-z*[*of K' v*], *assumption*, *erule* *exE*, *simp*,
cut-tac $n = N$ **in** *an-nat-pos*,
frule-tac $w1 = za$ **and** $x1 = 0$ **and** $y1 = an\ N$ **in**
amult-pos-mono-l[*THEN sym*], *simp*, *simp* *add:amult-0-r*)

apply (*frule-tac* $j = ant\ za * an\ N$ **and** $k = n\text{-val}\ K\ v\ (f\ n \pm_{K'} -_{a_{K'}} (f\ m))$
in *ale-trans*[*of 0*], *assumption*+)

apply (*frule-tac* $w1 = z$ **and** $x1 = 0$ **and** $y1 = n\text{-val}\ K\ v\ (f\ n \pm_{K'} -_{a_{K'}} (f\ m))$
in *amult-pos-mono-r*[*THEN sym*], *simp*, *simp* *add:amult-0-l*)

apply (*frule-tac* $i = Lv\ K'\ v' * an\ N$ **and** $j = n\text{-val}\ K\ v\ (f\ n \pm_{K'} -_{a_{K'}} (f\ m))$
and $k = v\ (f\ n \pm_{K'} -_{a_{K'}} (f\ m))$ **in** *ale-trans*, *assumption*+)

apply (*thin-tac* $f\ n \pm_{K'} -_{a_{K'}} (f\ m) \in vp\ K\ v\ (Vr\ K\ v)\ (Lv\ K'\ v') * (an\ N)$,
thin-tac $Lv\ K'\ v' * an\ N \leq n\text{-val}\ K\ v\ (f\ n \pm_{K'} -_{a_{K'}} (f\ m))$,
thin-tac $n\text{-val}\ K\ v\ (f\ n \pm_{K'} -_{a_{K'}} (f\ m)) \leq v\ (f\ n \pm_{K'} -_{a_{K'}} (f\ m))$)

apply (*simp* *add:v-completion-def*, (*erule* *conjE*)+)

apply (*thin-tac* $\forall x \in carrier\ K. v\ x = v'\ x$,
thin-tac $\forall x \in carrier\ K'. \exists f. Cauchy\ K\ v\ f \wedge \lim\ K'\ v'\ f\ x$)

apply (*frule* *subfield-sub*[*of K*],
frule-tac $c = f\ n \pm_{K'} -_{a_{K'}} (f\ m)$ **in**
subsetD[*of carrier K carrier K'*], *assumption*+)

apply (*simp* *add:Corps.n-val*[*THEN sym*, *of K' v*])

apply (*simp* *add:amult-commute*[*of - Lv K' v*])

apply (*frule* *Corps.Lv-pos*[*of K' v*], *assumption*,
frule *Corps.Lv-z*[*of K' v*], *assumption*+, *erule* *exE*, *simp*)

apply (*simp* *add:amult-pos-mono-l*)

apply (*rule-tac* $x = f\ n \pm_{K'} -_{a_{K'}} (f\ m)$ **and** $n = an\ N$ **in**
Corps.n-value-x-2[*of K' v*], *assumption*+)

apply (*cut-tac* $n = N$ **in** *an-nat-pos*)

apply (*frule-tac* $j = an\ N$ **and** $k = n\text{-val}\ K'\ v'\ (f\ n \pm_{K'} -_{a_{K'}} (f\ m))$ **in**
ale-trans[*of 0*], *assumption*+)

apply (*simp* *add:Corps.val-pos-n-val-pos*[*THEN sym*, *of K' v*])

apply (*simp* *add:Corps.val-pos-mem-Vr*) **apply** *assumption* **apply** *simp*
done

lemma *max-gtTr*:($n::nat$) < *max* (*Suc n*) (*Suc m*) $\wedge m$ < *max* (*Suc n*) (*Suc m*)
by (*simp* *add:max-def*)

lemma (**in** *Corps*) *completion-approx*: $\llbracket Corps\ K'; valuation\ K\ v; valuation\ K'\ v';$
Completion _{$v\ v'$} $K\ K'; x \in carrier\ (Vr\ K'\ v')$ $\implies$
 $\exists y \in carrier\ (Vr\ K\ v). (y \pm_{K'} -_{a_{K'}} x) \in (vp\ K'\ v')$

apply (*frule* *Corps.Vr-mem-f-mem* [*of K' v' x*], *assumption*+,

$\text{frule Corps.val-pos-mem-Vr}[\text{THEN sym, of } K' v' x], \text{assumption+},$
 $\text{simp add:v-completion-def, (erule conjE)+},$
 $\text{rotate-tac } -1, \text{drule-tac } x = x \text{ in bspec, assumption+},$
 $\text{erule exE, erule conjE})$
apply ($\text{unfold Cauchy-seq-def, frule conjunct1, fold Cauchy-seq-def}$)
apply ($\text{case-tac } x = \mathbf{0}_{K'},$
 $\text{simp, frule Corps.field-is-ring[of } K'],$
 $\text{frule Ring.ring-is-ag[of } K'],$
 $\text{subgoal-tac } \mathbf{0}_{K'} \in \text{carrier } (Vr K v),$
 $\text{subgoal-tac } (\mathbf{0}_{K'} \pm_{K'} -_a K' \mathbf{0}_{K'}) \in vp K' v', \text{blast},$
 $\text{frule aGroup.ag-inc-zero[of } K'], \text{simp add:aGroup.ag-r-inv1},$
 $\text{simp add:Corps.Vr-0-f-0[THEN sym, of } K' v'],$
 $\text{frule Corps.Vr-ring[of } K' v'], \text{assumption+},$
 $\text{frule Corps.vp-ideal[of } K' v'], \text{assumption+},$
 $\text{simp add:Ring.ideal-zero},$
 $\text{simp add:subfield-zero[THEN sym, of } K'],$
 $\text{cut-tac field-is-ring, frule Ring.ring-is-ag[of } K'],$
 $\text{frule aGroup.ag-inc-zero[of } K'],$
 $\text{simp add:Vr-0-f-0[THEN sym, of } v],$
 $\text{frule Vr-ring[of } v], \text{simp add:Ring.ring-zero})$
apply ($\text{frule-tac } f = f \text{ in Corps.limit-val[of } K' x - v'],$
 $\text{assumption+})$
apply ($\text{rule allI, rotate-tac } -2, \text{frule-tac } x = j \text{ in spec,}$
 $\text{frule subfield-sub[of } K'], \text{simp add:subsetD, assumption+})$
apply (erule exE)
apply ($\text{simp add:limit-def,}$
 $\text{frule Corps.Vr-ring[of } K' v'], \text{assumption+},$
 $\text{rotate-tac } 10,$
 $\text{drule-tac } x = \text{Suc } 0 \text{ in spec, erule exE},$
 $\text{rotate-tac } 1,$
 $\text{frule-tac } x = N \text{ and } y = M \text{ in two-inequalities, assumption+},$
 $\text{thin-tac } \forall n > N. v' (f n) = v' x,$
 $\text{thin-tac } \forall n > M. f n \pm_{K'} -_a K' x \in vp K' v' (Vr K' v') \text{ (an (Suc } 0))$)
apply ($\text{frule Corps.vp-ideal[of } K' v'], \text{assumption+},$
 $\text{simp add:Ring.r-apow-Suc[of } Vr K' v' vp K' v']$)
apply ($\text{drule-tac } x = N + M + 1 \text{ in spec, simp,}$
 $\text{drule-tac } x = N + M + 1 \text{ in spec, simp,}$
 $\text{erule conjE})$
apply ($\text{drule-tac } x = f (\text{Suc } (N + M)) \text{ in bspec, assumption+})$
apply simp
apply ($\text{cut-tac } x = f (\text{Suc } (N + M)) \text{ in val-pos-mem-Vr[of } v], \text{assumption+})$
apply simp apply blast
done

lemma (**in** Corps) $\text{res-v-completion-surj} : \llbracket \text{Corps } K'; \text{valuation } K v;$
 $\text{valuation } K' v'; \text{Completion}_{v v'} K K' \rrbracket \implies$
 $\text{surjec}^{(Vr K v), (qring (Vr K' v') (vp K' v'))}$
 $(\text{compos } (Vr K v) (pj (Vr K' v') (vp K' v')) (I_{(Vr K v)}))$
apply ($\text{frule Vr-ring[of } v],$

```

    frule Ring.ring-is-ag[of Vr K v],
    frule Corps.Vr-ring[of K' v'], assumption+,
    frule Ring.ring-is-ag[of Vr K' v'],
    frule Ring.ring-is-ag[of Vr K v])
apply (frule Corps.vp-ideal[of K' v'], assumption+,
    frule Ring.qring-ring[of Vr K' v' vp K' v'], assumption+)
apply (simp add:surjec-def)
apply (frule aHom-compos[of Vr K v Vr K' v'
    qring (Vr K' v') (vp K' v') I(Vr K v)
    pj (Vr K' v') (vp K' v')], assumption+, simp add:Ring.ring-is-ag)
apply (rule Vr-idmap-aHom, assumption+) apply (simp add:completion-subfield,
    simp add:v-completion-def) apply (
    frule pj-Hom[of Vr K' v' vp K' v'], assumption+) apply (
    simp add:rHom-def, simp)
apply (rule surj-to-test)
apply (simp add:aHom-def)
apply (rule ballI)
apply (thin-tac Ring (Vr K' v' /r vp K' v'),
    thin-tac compos (Vr K v) (pj (Vr K' v') (vp K' v')) (I(Vr K v)) ∈
    aHom (Vr K v) (Vr K' v' /r vp K' v'))
apply (simp add:Ring.qring-carrier)
apply (erule bexE)
apply (frule-tac x = a in completion-approx[of K' v v'],
    assumption+, erule bexE)
apply (subgoal-tac compos (Vr K v) (pj (Vr K' v')
    (vp K' v')) ((I(Vr K v))) y = b, blast)
apply (simp add:compos-def compose-def idmap-def)
apply (frule completion-Vring-sub[of K' v v'], assumption+)
apply (frule-tac c = y in subsetD[of carrier (Vr K v) carrier (Vr K' v')],
    assumption+)
apply (frule-tac x = y in pj-mem[of Vr K' v' vp K' v'], assumption+, simp,
    thin-tac pj (Vr K' v') (vp K' v') y = y ⊔(Vr K' v') (vp K' v'))
apply (rotate-tac -5, frule sym, thin-tac a ⊔(Vr K' v') (vp K' v') = b,
    simp)
apply (rule-tac b1 = y and a1 = a in Ring.ar-coset-same1[THEN sym,
    of Vr K' v' vp K' v'], assumption+)
apply (frule Ring.ring-is-ag[of Vr K' v'],
    frule-tac x = a in aGroup.ag-mOp-closed[of Vr K' v'],
    assumption+)
apply (simp add:Corps.Vr-mOp-f-mOp, simp add:Corps.Vr-pOp-f-pOp)
done

lemma (in Corps) res-v-completion-ker: [Corps K'; valuation K v;
    valuation K' v'; Completionv v' K K'] ⇒
    ker(Vr K v), (qring (Vr K' v') (vp K' v'))
    (compos (Vr K v) (pj (Vr K' v') (vp K' v')) (I(Vr K v))) = vp K v
apply (rule equalityI)
apply (rule subsetI)

```

```

apply (simp add:ker-def, (erule conjE)+)
apply (frule Corps.Vr-ring[of  $K' v'$ ], assumption+,
      frule Corps.vp-ideal[of  $K' v'$ ], assumption+,
      frule Ring.qring-ring[of  $Vr K' v' vp K' v'$ ], assumption+,
      simp add:Ring.qring-zero)
apply (simp add:compos-def, simp add:compose-def, simp add:idmap-def)
apply (frule completion-Vring-sub[of  $K' v v'$ ], assumption+)
apply (frule-tac  $c = x$  in subsetD[of carrier ( $Vr K v$ ) carrier ( $Vr K' v'$ )],
      assumption+)
apply (simp add:pj-mem)
apply (frule-tac  $a = x$  in Ring.qring-zero-1[of  $Vr K' v' - vp K' v'$ ],
      assumption+)
apply (subst vp-mem-val-poss[of  $v$ ], assumption+)
apply (simp add:Vr-mem-f-mem)
apply (frule-tac  $x = x$  in Corps.vp-mem-val-poss[of  $K' v'$ ],
      assumption+, simp add:Corps.Vr-mem-f-mem, simp)
apply (frule-tac  $x = x$  in Vr-mem-f-mem[of  $v$ ], assumption+)
apply (frule-tac  $x = x$  in completion-val-eq[of  $K' v v'$ ],
      assumption+, simp)
apply (rule subsetI)
apply (simp add:ker-def)
apply (frule Vr-ring[of  $v$ ])
apply (frule vp-ideal[of  $v$ ])
apply (frule-tac  $h = x$  in Ring.ideal-subset[of  $Vr K v vp K v$ ],
      assumption+, simp)
apply (frule Corps.Vr-ring[of  $K' v'$ ], assumption+,
      frule Corps.vp-ideal[of  $K' v'$ ], assumption+,
      simp add:Ring.qring-zero)
apply (simp add:compos-def compose-def idmap-def)
apply (frule completion-Vring-sub[of  $K' v v'$ ],
      assumption+, frule-tac  $c = x$  in subsetD[of carrier ( $Vr K v$ )
      carrier ( $Vr K' v'$ )], assumption+)
apply (simp add:pj-mem)
apply (frule completion-vpr-sub[of  $K' v v'$ ], assumption+,
      frule-tac  $c = x$  in subsetD[of  $vp K v vp K' v'$ ], assumption+)
apply (simp add:Ring.ar-coset-same4[of  $Vr K' v' vp K' v'$ ])
done

```

```

lemma (in Corps) completion-res-qring-isom: [Corps  $K'$ ; valuation  $K v$ ;
  valuation  $K' v'$ ; Completion $v v'$   $K K'$ ]  $\implies$ 
  r-isom (( $Vr K v$ ) / $r$  ( $vp K v$ )) (( $Vr K' v'$ ) / $r$  ( $vp K' v'$ ))
apply (subst r-isom-def)
apply (frule res-v-completion-surj[of  $K' v v'$ ], assumption+)
apply (frule Vr-ring[of  $v$ ],
      frule Corps.Vr-ring[of  $K' v'$ ], assumption+,
      frule Corps.vp-ideal[of  $K' v'$ ], assumption+,
      frule Ring.qring-ring[of  $Vr K' v' vp K' v'$ ], assumption+)
apply (frule rHom-compos[of  $Vr K v Vr K' v' (Vr K' v' /_r vp K' v')$ 
  ( $I_{(Vr K v)}$ ) pj ( $Vr K' v')$  ( $vp K' v'$ )], assumption+)

```

apply (*simp add:completion-idmap-rHom*)
apply (*simp add:pj-Hom*)
apply (*frule surjec-ind-bijec [of Vr K v (Vr K' v' /_r vp K' v')*
compos (Vr K v) (pj (Vr K' v') (vp K' v')) (I(Vr K v))], assumption+)
apply (*frule ind-hom-rhom[of Vr K v (Vr K' v' /_r vp K' v')*
compos (Vr K v) (pj (Vr K' v') (vp K' v')) (I(Vr K v))], assumption+)
apply (*simp add:res-v-completion-ker*) **apply** *blast*
done

expansion of x in a complete field K , with normal valuation v . Here we suppose t is an element of K satisfying the equation $v\ t = 1$.

definition

$Kx :: [-, 'b \Rightarrow ant, 'b] \Rightarrow 'b$ set **where**
 $Kx\ K\ v\ x = \{y. \exists k \in carrier\ (Vr\ K\ v). y = x \cdot_r K\ k\}$

primrec

$partial-sum :: [('b, 'm)\ Ring-scheme, 'b, 'b \Rightarrow ant, 'b]$
 $\Rightarrow nat \Rightarrow 'b$
 $(\langle (5psum\ -\ -\ -\ -) \rangle [96,96,96,96,97]96)$

where

$psum-0: psum\ K\ x\ v\ t\ 0 = (csrp-fn\ (Vr\ K\ v)\ (vp\ K\ v)$
 $(pj\ (Vr\ K\ v)\ (vp\ K\ v)\ (x \cdot_r K\ t_K^{-(tna\ (v\ x))}))) \cdot_r K\ (t_K^{(tna\ (v\ x))})$
 $| psum-Suc: psum\ K\ x\ v\ t\ (Suc\ n) = (psum\ K\ x\ v\ t\ n) \pm_K$
 $((csrp-fn\ (Vr\ K\ v)\ (vp\ K\ v)\ (pj\ (Vr\ K\ v)\ (vp\ K\ v)$
 $((x \pm_K -_a K\ (psum\ K\ x\ v\ t\ n)) \cdot_r K\ (t_K^{-(tna\ (v\ x) + int\ (Suc\ n))})))) \cdot_r K$
 $(t_K^{(tna\ (v\ x) + int\ (Suc\ n))}))$

definition

$expand-coeff :: [-, 'b \Rightarrow ant, 'b, 'b]$
 $\Rightarrow nat \Rightarrow 'b$
 $(\langle (5ecf\ -\ -\ -\ -) \rangle [96,96,96,96,97]96)$ **where**
 $ecf_{K\ v\ t\ x}\ n = (if\ n = 0\ then\ csrp-fn\ (Vr\ K\ v)\ (vp\ K\ v)$
 $(pj\ (Vr\ K\ v)\ (vp\ K\ v)\ (x \cdot_r K\ t_K^{-(tna\ (v\ x))})))$
 $else\ csrp-fn\ (Vr\ K\ v)\ (vp\ K\ v)\ (pj\ (Vr\ K\ v)$
 $(vp\ K\ v)\ ((x \pm_K -_a K\ (psum\ K\ x\ v\ t\ (n - 1))) \cdot_r K\ (t_K^{-(tna\ (v\ x) + int\ n)}))))$

definition

$expand-term :: [-, 'b \Rightarrow ant, 'b, 'b]$
 $\Rightarrow nat \Rightarrow 'b$
 $(\langle (5etm\ -\ -\ -\ -) \rangle [96,96,96,96,97]96)$ **where**

$etm_{K\ v\ t\ x}\ n = (ecf_{K\ v\ t\ x}\ n) \cdot_r K\ (t_K^{(tna\ (v\ x) + int\ n)})$

```

lemma (in Corps) Kxa-val-ge: [valuation K v; t ∈ carrier K; v t = 1]
  ⇒ Kxa K v (t_K^j) = {x. x ∈ carrier K ∧ (ant j) ≤ (v x)}
apply (frule valI-neg-0[of v t], assumption+)
apply (cut-tac field-is-ring)
apply (rule equalityI)
apply (rule subsetI,
  simp add:Kxa-def, erule bexE,
  frule-tac x = k in Vr-mem-f-mem[of v], assumption+,
  frule npowf-mem[of t j], simp,
  simp add:Ring.ring-tOp-closed)
apply (simp add:val-t2p) apply (simp add:val-exp[THEN sym])
apply (simp add:val-pos-mem-Vr[THEN sym, of v])
apply (frule-tac x = 0 and y = v k in aadd-le-mono[of - - ant j])
apply (simp add:aadd-0-l aadd-commute)
apply (rule subsetI, simp, erule conjE)
apply (simp add:Kxa-def)
apply (case-tac x = 0_K)
apply (frule Vr-ring[of v],
  frule Ring.ring-zero[of Vr K v],
  simp add:Vr-0-f-0,
  frule-tac x1 = t_K^j in Ring.ring-times-x-0[THEN sym, of K],
  simp add:npowf-mem, blast)
apply (frule val-exp[of v t j], assumption+, simp)
apply (frule field-potent-nonzero1[of t j],
  frule npowf-mem[of t j], assumption+)
apply (frule-tac y = v x in ale-diff-pos[of v (t_K^j)],
  simp add:diff-ant-def)
apply (cut-tac npowf-mem[of t j])
defer
apply assumption apply simp
apply (frule value-of-inv[THEN sym, of v t_K^j], assumption+)

apply (cut-tac invf-closed1[of t_K^j], simp, erule conjE)
apply (frule-tac x1 = x in val-t2p[THEN sym, of v - (t_K^j)^-K],
  assumption+, simp)
apply (frule-tac x = (t_K^j)^-K and y = x in Ring.ring-tOp-closed[of K],
  assumption+)
apply (simp add:Ring.ring-tOp-commute[of K - (t_K^j)^-K])
apply (frule-tac x = ((t_K^j)^-K) ·r x in val-pos-mem-Vr[of v], assumption+,
  simp)
apply (frule-tac z = x in Ring.ring-tOp-assoc[of K t_K^j (t_K^j)^-K],
  assumption+)
apply (simp add:Ring.ring-tOp-commute[of K t_K^j (t_K^j)^-K] linvf,
  simp add:Ring.ring-l-one, blast)
apply simp
done

```

```

lemma (in Corps) Kxa-pow-vpr: [valuation K v; t ∈ carrier K; v t = 1;

```

$(0::int) \leq j \implies Kxa \ K \ v \ (t_K^j) = (vp \ K \ v)(Vr \ K \ v) \ (ant \ j)$
apply (frule val-surj-n-val[of v], blast)
apply (simp add:Kxa-val-ge)
apply (rule equalityI)
apply (rule subsetI, simp, erule conjE)
apply (rule-tac $x = x$ **in** n-value-x-2[of v - (ant j)],
 assumption+)
apply (cut-tac ale-zle[THEN sym, of 0 j])
apply (frule-tac $a = 0 \leq j$ **and** $b = ant \ 0 \leq ant \ j$ **in** a-b-exchange,
 assumption+)
apply (thin-tac $(0 \leq j) = (ant \ 0 \leq ant \ j)$, simp add:ant-0)
apply (frule-tac $k = v \ x$ **in** ale-trans[of 0 ant j], assumption+)
apply (simp add:val-pos-mem-Vr) **apply** simp
apply (simp only:ale-zle[THEN sym, of 0 j], simp add:ant-0)
apply (rule subsetI)
apply simp
apply (frule-tac $x = x$ **in** mem-vp-apow-mem-Vr[of v ant j])
apply (simp only:ale-zle[THEN sym, of 0 j], simp add:ant-0)
apply assumption
apply (simp add:Vr-mem-f-mem[of v])
apply (frule-tac $x = x$ **in** n-value-x-1[of v ant j -])
apply (simp only:ale-zle[THEN sym, of 0 j], simp add:ant-0)
apply assumption **apply** simp
done

lemma (in Corps) field-distribTr: $\llbracket a \in carrier \ K; \ b \in carrier \ K; \$
 $x \in carrier \ K; \ x \neq \mathbf{0} \rrbracket \implies a \pm (-_a \ (b \cdot_r \ x)) = (a \cdot_r \ (x^K) \pm (-_a \ b)) \cdot_r \ x$
apply (cut-tac field-is-ring,
 cut-tac invf-closed1[of x], simp, erule conjE) **apply** (
 simp add:Ring.ring-inv1-1[of K b x],
 frule Ring.ring-is-ag[of K],
 frule aGroup.ag-mOp-closed[of K b], assumption+)
apply (frule Ring.ring-tOp-closed[of K a x^K], assumption+,
 simp add:Ring.ring-distrib2, simp add:Ring.ring-tOp-assoc,
 simp add:linvf,
 simp add:Ring.ring-r-one)
apply simp
done

lemma a0-le-1[simp]: $(0::ant) \leq 1$
by (cut-tac a0-less-1, simp add:aless-imp-le)

lemma (in Corps) vp-mem-times-t: $\llbracket valuation \ K \ v; \ t \in carrier \ K; \ t \neq \mathbf{0}; \$
 $v \ t = 1; \ x \in vp \ K \ v \rrbracket \implies \exists a \in carrier \ (Vr \ K \ v). \ x = a \cdot_r \ t$
apply (frule Vr-ring[of v],
 frule vp-ideal[of v])
apply (frule val-surj-n-val[of v], blast)
apply (frule vp-mem-val-poss[of v x],
 frule-tac $h = x$ **in** Ring.ideal-subset[of Vr K v vp K v],
 assumption+)

```

      assumption+, simp add: Vr-mem-f-mem, simp)
apply (frule gt-a0-ge-1[of v x])
apply (subgoal-tac  $v \leq v$ )
apply (thin-tac  $1 \leq v$ )
apply (frule val-Rxa-gt-a-1[of v t x])
apply (subst val-pos-mem-Vr[THEN sym, of v t], assumption+)
apply simp
apply (simp add:vp-mem-Vr-mem) apply assumption+
apply (simp add:Rxa-def, erule bexE)
apply (cut-tac a0-less-1)
apply (subgoal-tac  $0 \leq v$ )
apply (frule val-pos-mem-Vr[of v t], assumption+)
apply (simp, simp add: Vr-tOp-f-tOp, blast, simp+)
done

lemma (in Corps) psum-diff-mem-Kxa:  $\llbracket \text{valuation } K \ v; \ t \in \text{carrier } K; \$ 
 $v \ t = 1; \ x \in \text{carrier } K; \ x \neq \mathbf{0} \rrbracket \implies$ 
 $(\text{psum } K \ x \ v \ t \ n) \in \text{carrier } K \wedge$ 
 $(x \pm (-_a (\text{psum } K \ x \ v \ t \ n))) \in Kxa \ K \ v \ (t_K^{((tna \ (v \ x)) + (1 + \text{int } n))})$ 
apply (frule val1-neg-0[of v t], assumption+)
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K],
      frule Vr-ring[of v], frule vp-ideal[of v])
apply (induct-tac  $n$ )
apply (subgoal-tac  $x \cdot_r (t_K^{- \ tna \ (v \ x)}) \in \text{carrier } (Vr \ K \ v),$ 
      rule conjI, simp, rule Ring.ring-tOp-closed[of K], assumption+,
      frule Ring.csrp-fn-mem[of Vr K v vp K v
      pj (Vr K v) (vp K v) ( $x \cdot_r (t_K^{- \ tna \ (v \ x)})$ )],
      assumption+,
      simp add: pj-mem Ring.qring-carrier, blast,
      simp add: Vr-mem-f-mem, simp add: npowf-mem)
apply (simp,
      frule npowf-mem[of t tna (v x)], assumption+,
      frule field-potent-nonzero1[of t tna (v x)], assumption+,
      subst field-distribTr[of x - t_K^{(tna (v x))}], assumption+,
      frule Ring.csrp-fn-mem[of Vr K v vp K v
      pj (Vr K v) (vp K v) ( $x \cdot_r (t_K^{- \ tna \ (v \ x)})$ )],
      assumption+,
      simp add: pj-mem Ring.qring-carrier, blast, simp add: Vr-mem-f-mem,
      simp add: npowf-mem, assumption)
apply (frule Ring.csrp-diff-in-vpr[of Vr K v vp K v
       $x \cdot_r ((t_K^{(tna \ (v \ x))})^{-K})$ ], assumption+)
apply (simp add: npowf-minus)

apply (simp add: npowf-minus)
apply (frule pj-Hom[of Vr K v vp K v], assumption+)
apply (frule rHom-mem[of pj (Vr K v) (vp K v) Vr K v Vr K v /_r vp K v
       $x \cdot_r (t_K^{- \ tna \ (v \ x)})$ ], assumption+)
apply (frule Ring.csrp-fn-mem[of Vr K v vp K v

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      pj (Vr K v) (vp K v) (x ·r (tK- tna (v x))), assumption+)
apply (frule Ring.ring-is-ag[of Vr K v],
      frule-tac x = csrp-fn (Vr K v) (vp K v) (pj (Vr K v) (vp K v)
      (x ·r (tK- tna (v x)))) in aGroup.ag-mOp-closed[of Vr K v], assumption+)
apply (simp add: Vr-pOp-f-pOp) apply (simp add: Vr-mOp-f-mOp)
apply (frule-tac x = x ·r (tK- tna (v x)) ± -a (csrp-fn (Vr K v) (vp K v)
      (pj (Vr K v) (vp K v) (x ·r (tK- tna (v x))))) in
      vp-mem-times-t[of v t], assumption+, erule bexE, simp)
apply (frule-tac x = a in Vr-mem-f-mem[of v], assumption+)
apply (simp add: Ring.ring-tOp-assoc[of K])
apply (simp add: npowf-exp-1-add[THEN sym, of t tna (v x)])
apply (simp add: add commute)
apply (simp add: Kxa-def)
apply (frule npowf-mem[of t 1 + tna (v x)], assumption+)
apply (simp add: Ring.ring-tOp-commute[of K - tK(1 + tna (v x))])
apply blast
apply (frule npowf-mem[of t - tna (v x)], assumption+)
apply (frule Ring.ring-tOp-closed[of K x tK- tna (v x)], assumption+)
apply (subst val-pos-mem-Vr[THEN sym, of v], assumption+)
apply (simp add: val-t2p) apply (simp add: val-exp[THEN sym, of v t])
apply (simp add: aminus[THEN sym])
apply (frule value-in-aug-inf[of v x], assumption+,
      simp add: aug-inf-def)
apply (frule val-nonzero-noninf[of v x], assumption+,
      simp add: ant-tna)
apply (simp add: aadd-minus-r)

apply (erule conjE)
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (rule conjI)
apply simp
apply (rule aGroup.ag-pOp-closed, assumption+)
apply (rule Ring.ring-tOp-closed[of K], assumption+)
apply (simp add: Kxa-def, erule bexE, simp)
apply (subst Ring.ring-tOp-commute, assumption+)
apply (rule npowf-mem, assumption+) apply (simp add: Vr-mem-f-mem)
apply (frule-tac n = tna (v x) + (1 + int n) in npowf-mem[of t ],
      assumption,
      frule-tac n = - tna (v x) + (-1 - int n) in npowf-mem[of t ],
      assumption,
      frule-tac x = k in Vr-mem-f-mem[of v], assumption+)
apply (simp add: Ring.ring-tOp-assoc npowf-exp-add[THEN sym, of t])
apply (simp add: field-mpowf-exp-zero)
apply (simp add: Ring.ring-r-one)

apply (frule pj-Hom[of Vr K v vp K v], assumption+)
apply (frule-tac a = k in rHom-mem[of pj (Vr K v) (vp K v) Vr K v
      Vr K v /r vp K v], assumption+)

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apply (frule-tac  $x = pj (Vr K v) (vp K v) k$  in Ring.csrp-fn-mem[of  $Vr K v$ 
   $vp K v$ ], assumption+)
apply (simp add: Vr-mem-f-mem)
apply (rule npowf-mem, assumption+)

apply (simp add: Kxa-def) apply (erule bexE, simp)
apply (frule-tac  $x = k$  in Vr-mem-f-mem[of  $v$ ], assumption+)
apply (frule-tac  $n = tna (v x) + (1 + int n)$  in npowf-mem[of  $t$ ],
  assumption+)
apply (frule-tac  $n = - tna (v x) + (- 1 - int n)$  in npowf-mem[of  $t$ ],
  assumption+)
apply (frule-tac  $x = t_K^{(tna (v x) + (1 + int n))}$  and  $y = k$  in
  Ring.ring-tOp-commute[of  $K$ ], assumption+) apply simp
apply (simp add: Ring.ring-tOp-assoc, simp add: npowf-exp-add[THEN sym])
  apply (simp add: field-mpowf-exp-zero)
  apply (simp add: Ring.ring-r-one)
  apply (thin-tac ( $t_K^{(tna (v x) + (1 + int n))}$ )  $\cdot_r k =$ 
     $k \cdot_r (t_K^{(tna (v x) + (1 + int n))})$ )
apply (frule pj-Hom[of  $Vr K v vp K v$ ], assumption+)
apply (frule-tac  $a = k$  in rHom-mem[of  $pj (Vr K v) (vp K v) Vr K v$ 
   $Vr K v /_r vp K v$ ], assumption+)
apply (frule-tac  $x = pj (Vr K v) (vp K v) k$  in Ring.csrp-fn-mem[of  $Vr K v$ 
   $vp K v$ ], assumption+)
apply (frule-tac  $x = csrp-fn (Vr K v) (vp K v) (pj (Vr K v) (vp K v) k)$  in
  Vr-mem-f-mem[of  $v$ ], assumption+)
apply (frule-tac  $x = csrp-fn (Vr K v) (vp K v) (pj (Vr K v) (vp K v) k)$  and
   $y = t_K^{(tna (v x) + (1 + int n))}$  in Ring.ring-tOp-closed[of  $K$ ], assumption+)
apply (simp add: aGroup.ag-p-inv[of  $K$ ])
apply (frule-tac  $x = psum K x v t n$  in aGroup.ag-mOp-closed[of  $K$ ],
  assumption+)
apply (frule-tac  $x = (csrp-fn (Vr K v) (vp K v) (pj (Vr K v) (vp K v) k)) \cdot_r$ 
  ( $t_K^{(tna (v x) + (1 + int n))}$ ) in aGroup.ag-mOp-closed[of  $K$ ], assumption+)
apply (subst aGroup.ag-pOp-assoc[THEN sym], assumption+)
apply simp
apply (simp add: Ring.ring-inv1-1)
apply (subst Ring.ring-distrib2[THEN sym, of  $K$ ], assumption+)
apply (rule aGroup.ag-mOp-closed, assumption+)
apply (frule-tac  $x = k$  in Ring.csrp-diff-in-vpr[of  $Vr K v vp K v$ 
  ], assumption+)
apply (frule Ring.ring-is-ag[of  $Vr K v$ ])
apply (frule-tac  $x = csrp-fn (Vr K v) (vp K v) (pj (Vr K v) (vp K v) k)$  in
  aGroup.ag-mOp-closed[of  $Vr K v$ ], assumption+)
apply (simp add: Vr-pOp-f-pOp) apply (simp add: Vr-mOp-f-mOp)
apply (frule-tac  $x = k \pm -_a (csrp-fn (Vr K v) (vp K v) (pj (Vr K v) (vp K v) k))$  in
  vp-mem-times-t[of  $v t$ ], assumption+, erule bexE, simp)
apply (frule-tac  $x = a$  in Vr-mem-f-mem[of  $v$ ], assumption+)
apply (simp add: Ring.ring-tOp-assoc[of  $K$ ])
apply (simp add: npowf-exp-1-add[THEN sym, of  $t$ ])

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apply (simp add:add.commute[of 2])
apply (simp add:add.assoc)
apply (subst Ring.ring-tOp-commute, assumption+)
apply (rule npowf-mem, assumption+) apply blast
done

lemma Suc-diff-int:  $0 < n \implies \text{int } (n - \text{Suc } 0) = \text{int } n - 1$ 
by (cut-tac of-nat-Suc[of n - Suc 0], simp)

lemma (in Corps) ecf-mem:  $\llbracket \text{valuation } K \ v; t \in \text{carrier } K; v \ t = 1; \\ x \in \text{carrier } K; x \neq \mathbf{0} \rrbracket \implies \text{ecf}_K \ v \ t \ x \ n \in \text{carrier } K$ 
apply (frule val1-neg-0[of v t], assumption+)
apply (cut-tac field-is-ring,
  frule Vr-ring[of v], frule vp-ideal[of v])
apply (case-tac n = 0)
apply (simp add:expand-coeff-def)
apply (rule Vr-mem-f-mem[of v], assumption+)
apply (rule Ring.csrp-fn-mem, assumption+)
apply (subgoal-tac  $x \cdot_r (t_K^{- \text{tna } (v \ x)}) \in \text{carrier } (Vr \ K \ v)$ )
apply (simp add:pj-mem Ring.qring-carrier, blast)
apply (frule npowf-mem[of t - tna (v x)], assumption+,
  subst val-pos-mem-Vr[THEN sym, of  $v \ x \cdot_r (t_K^{(- \text{tna } (v \ x))})$ ],
  assumption+,
  rule Ring.ring-tOp-closed, assumption+)
apply (simp add:val-t2p,
  simp add:val-exp[THEN sym, of  $v \ t - \text{tna } (v \ x)$ ])
apply (frule value-in-aug-inf[of v x], assumption+,
  simp add:aug-inf-def)
apply (frule val-nonzero-noninf[of v x], assumption+)
apply (simp add:aminus[THEN sym], simp add:ant-tna)
apply (simp add:aadd-minus-r)

apply (simp add:expand-coeff-def)
apply (frule psum-diff-mem-Kxa[of v t x n - 1],
  assumption+, erule conjE)
apply (simp add:Kxa-def, erule bexE,
  frule-tac  $x = k$  in Vr-mem-f-mem[of v], assumption+,
  frule npowf-mem[of t tna (v x) + (1 + int (n - Suc 0))],
  assumption+,
  frule npowf-mem[of t - tna (v x) - int n], assumption+)
apply simp
apply (simp add:Ring.ring-tOp-commute[of  $K \ t_K^{(\text{tna } (v \ x) + (\text{int } n))}$ ])
apply (simp add:Ring.ring-tOp-assoc, simp add:npowf-exp-add[THEN sym])

apply (simp add:npowf-def, simp add:Ring.ring-r-one)
apply (rule Vr-mem-f-mem, assumption+)
apply (rule Ring.csrp-fn-mem, assumption+)
apply (simp add:pj-mem Ring.qring-carrier, blast)
done

```

lemma (in Corps) etm-mem: $\llbracket \text{valuation } K \ v; t \in \text{carrier } K; v \ t = 1; x \in \text{carrier } K; x \neq \mathbf{0} \rrbracket \implies \text{etm}_{K \ v \ t \ x} \ n \in \text{carrier } K$
apply (frule val1-neg-0[of v t], assumption+)
apply (simp add:expand-term-def)
apply (cut-tac field-is-ring,
 rule Ring.ring-tOp-closed[of K], assumption)
apply (simp add:ecf-mem)
apply (simp add:npowf-mem)
done

lemma (in Corps) psum-sum-etm: $\llbracket \text{valuation } K \ v; t \in \text{carrier } K; v \ t = 1; x \in \text{carrier } K; x \neq \mathbf{0} \rrbracket \implies$
 $\text{psum}_{K \ x \ v \ t} \ n = \text{nsun } K \ (\lambda j. (\text{ecf}_{K \ v \ t \ x} \ j) \cdot_r (t_K^{(\text{tna } (v \ x) + (\text{int } j))})) \ n$
apply (frule val1-neg-0[of v t], assumption+)
apply (induct-tac n)
apply (simp add:expand-coeff-def)
apply (rotate-tac -1, drule sym)
apply simp
apply (thin-tac $\Sigma_e \ K \ (\lambda j. \text{ecf}_{K \ v \ t \ x} \ j \cdot_r t_K^{(\text{tna } (v \ x) + (\text{int } j))}) \ n =$
 $\text{psum}_{K \ x \ v \ t} \ n$)
apply (simp add:expand-coeff-def)
done

lemma zabs-pos: $0 \leq (\text{abs } (z::\text{int}))$
by (simp add:zabs-def)

lemma abs-p-self-pos: $0 \leq z + (\text{abs } (z::\text{int}))$
by (simp add:zabs-def)

lemma zadd-right-mono: $(i::\text{int}) \leq j \implies i + k \leq j + k$
by simp

theorem (in Corps) expansion-thm: $\llbracket \text{valuation } K \ v; t \in \text{carrier } K; v \ t = 1; x \in \text{carrier } K; x \neq \mathbf{0} \rrbracket \implies \lim_{K \ v} (\text{partial-sum } K \ x \ v \ t) \ x$
apply (cut-tac field-is-ring, frule Ring.ring-is-ag[of K])
apply (simp add:limit-def)
apply (rule allI)
apply (subgoal-tac $\forall n. (N + \text{na } (\text{Abs } (v \ x))) < n \longrightarrow$
 $\text{psum}_{K \ x \ v \ t} \ n \pm -_a \ x \in \text{vp } K \ v^{(\text{Vr } K \ v)} (\text{an } N)$)
apply blast
apply (rule allI, rule impI)
apply (frule-tac $n = n$ in psum-diff-mem-Kxa[of v t x],
 assumption+, erule conjE)
apply (frule-tac $j = \text{tna } (v \ x) + (1 + \text{int } n)$ in Kxa-val-ge[of v t],
 assumption+)
apply simp
apply (thin-tac $Kxa \ K \ v \ (t_K^{(\text{tna } (v \ x) + (1 + \text{int } n))}) =$
 $\{xa \in \text{carrier } K. \text{ant } (\text{tna } (v \ x) + (1 + \text{int } n)) \leq v \ xa\}$)

```

apply (erule conjE)
apply (simp add:a-zpz[THEN sym])

apply (subgoal-tac (an N) ≤ v (psum K x v t n ± -a x))

apply (frule value-in-aug-inf[of v x], assumption+,
  simp add:aug-inf-def)
apply (frule val-nonzero-noninf[of v x], assumption+)
apply (simp add:ant-tna)
apply (frule val-surj-n-val[of v], blast)
apply (rule-tac x = psum K x v t n ± -a x and n = an N in
  n-value-x-2[of v], assumption+)
apply (subst val-pos-mem-Vr[THEN sym, of v], assumption+)
apply (rule aGroup.ag-pOp-closed[of K], assumption+)
apply (simp add:aGroup.ag-mOp-closed)

apply (cut-tac n = N in an-nat-pos)
apply (rule-tac i = 0 and j = an N and k = v (psum K x v t n ± -a x) in
  ale-trans, assumption+)
apply simp
apply simp

apply (frule-tac x1 = x ± (-a psum K x v t n) in val-minus-eq[THEN sym,
  of v], assumption+, simp,
  thin-tac v ( x ± -a psum K x v t n) = v (-a ( x ± -a psum K x v t n)))
apply (frule-tac x = psum K x v t n in aGroup.ag-mOp-closed[of K],
  assumption+, simp add:aGroup.ag-p-inv, simp add:aGroup.ag-inv-inv)
apply (frule aGroup.ag-mOp-closed[of K x], assumption+)
apply (simp add:aGroup.ag-pOp-commute[of K -a x])

apply (cut-tac Abs-pos[of v x])
apply (frule val-nonzero-z[of v x], assumption+, erule exE, simp)
apply (simp add:na-def)
apply (cut-tac aneg-less[THEN sym, of 0 Abs (v x)], simp)
apply (frule val-nonzero-noninf[of v x], assumption+)
apply (simp add:tna-ant)
apply (simp only:ant-1[THEN sym], simp del:ant-1 add:a-zpz)
apply (simp add:add.assoc[THEN sym])
apply (cut-tac m1 = N + nat |z| and n1 = n in of-nat-less-iff [where ?'a =
  int, THEN sym], simp)
apply (frule-tac a = int N + abs z and b = int n and c = z + 1 in
  add-strict-right-mono, simp only:add commute)
apply (simp only:add.assoc[of - 1])
apply (simp only:add.assoc[THEN sym, of 1])
apply (simp only:add commute[of 1])
apply (simp only:add.assoc[of - 1])
apply (cut-tac ?a1 = z and ?b1 = abs z and ?c1 = 1 + int N in
  add.assoc[THEN sym])
apply (thin-tac |z| + int N < int n)

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apply (*frule-tac* $a = z + (|z| + (1 + \text{int } N))$ **and** $b = z + |z| + (1 + \text{int } N)$ **and**
 $c = \text{int } n + (z + 1)$ **in** *ineq-conv1, assumption+*)
apply (*thin-tac* $z + (|z| + (1 + \text{int } N)) < \text{int } n + (z + 1)$,
 $\text{thin-tac } z + (|z| + (1 + \text{int } N)) = z + |z| + (1 + \text{int } N)$,
 $\text{thin-tac } N + \text{nat } |z| < n$)
apply (*cut-tac* $z = z$ **in** *abs-p-self-pos*)
apply (*frule-tac* $i = 0$ **and** $j = z + \text{abs } z$ **and** $k = 1 + \text{int } N$ **in**
 zadd-right-mono)
apply (*simp only:add-0*)
apply (*frule-tac* $i = 1 + \text{int } N$ **and** $j = z + |z| + (1 + \text{int } N)$ **and**
 $k = \text{int } n + (z + 1)$ **in** *zle-zless-trans, assumption+*)
apply (*thin-tac* $z + |z| + (1 + \text{int } N) < \text{int } n + (z + 1)$,
 $\text{thin-tac } 0 \leq z + |z|$,
 $\text{thin-tac } 1 + \text{int } N \leq z + |z| + (1 + \text{int } N)$)
apply (*cut-tac* $m1 = 1 + \text{int } N$ **and** $n1 = \text{int } n + (z + 1)$ **in**
 $\text{aless-zless}[THEN \text{sym}], \text{simp}$)

apply (*frule-tac* $x = \text{ant } (1 + \text{int } N)$ **and** $y = \text{ant } (\text{int } n + (z + 1))$
and $z = v (\text{psum } K \ x \ v \ t \ n \pm -_a \ x)$ **in** *aless-le-trans, assumption+*)
apply (*thin-tac* $\text{ant } (1 + \text{int } N) < \text{ant } (\text{int } n + (z + 1))$)

apply (*simp add:a-zpz[THEN sym]*)
apply (*frule-tac* $x = 1 + \text{ant } (\text{int } N)$ **and** $y = v (\text{psum } K \ x \ v \ t \ n \pm -_a \ x)$ **in**
 $\text{aless-imp-le}, \text{thin-tac } 1 + \text{ant } (\text{int } N) < v (\text{psum } K \ x \ v \ t \ n \pm -_a \ x)$)
apply (*cut-tac a0-less-1, frule aless-imp-le[of 0 1]*)
apply (*frule-tac* $b = \text{ant } (\text{int } N)$ **in** *aadd-pos-le[of 1]*)
apply (*subst an-def*)
apply (*rule-tac* $i = \text{ant } (\text{int } N)$ **and** $j = 1 + \text{ant } (\text{int } N)$ **and**
 $k = v (\text{psum } K \ x \ v \ t \ n \pm -_a \ x)$ **in** *ale-trans, assumption+*)
done

2.9.1 Hensel's theorem

definition

$\text{pol-Cauchy-seq} :: [(\text{'b}, \text{'m}) \text{ Ring-scheme}, \text{'b}, -, \text{'b} \Rightarrow \text{ant},$
 $\text{nat} \Rightarrow \text{'b}] \Rightarrow \text{bool} \ (\langle (5\text{PCauchy} \ \dots) \rangle [90,90,90,90,91]90) \text{ where}$
 $\text{PCauchy}_R \ X \ K \ v \ F \longleftrightarrow (\forall n. (F \ n) \in \text{carrier } R) \wedge$
 $(\exists d. (\forall n. \text{deg } R \ (Vr \ K \ v) \ X \ (F \ n) \leq (\text{an } d))) \wedge$
 $(\forall N. \exists M. (\forall n \ m. M < n \wedge M < m \longrightarrow$
 $P\text{-mod } R \ (Vr \ K \ v) \ X \ ((vp \ K \ v)(Vr \ K \ v) \ (\text{an } N)) \ (F \ n \pm_R -_a R \ (F \ m))))$

definition

$\text{pol-limit} :: [(\text{'b}, \text{'m}) \text{ Ring-scheme}, \text{'b}, -, \text{'b} \Rightarrow \text{ant},$
 $\text{nat} \Rightarrow \text{'b}, \text{'b}] \Rightarrow \text{bool}$
 $(\langle (6\text{Plimit} \ \dots) \rangle [90,90,90,90,90,91]90) \text{ where}$
 $\text{Plimit}_R \ X \ K \ v \ F \ p \longleftrightarrow (\forall n. (F \ n) \in \text{carrier } R) \wedge$
 $(\forall N. \exists M. (\forall m. M < m \longrightarrow$
 $P\text{-mod } R \ (Vr \ K \ v) \ X \ ((vp \ K \ v)(Vr \ K \ v) \ (\text{an } N)) \ ((F \ m) \pm_R -_a R \ p)))$

$$Pseq_{\ell} :: [(('b, 'm) \text{ Ring-scheme}, 'b, -, 'b \Rightarrow \text{ant}, \text{nat}, \\ \text{nat} \Rightarrow 'b) \Rightarrow \text{nat} \Rightarrow 'b \\ (\langle (6Pseq_{\ell} \text{ } \dots \text{ } \rangle) [90, 90, 90, 90, 90, 91] 90) \text{ where } \\ Pseq_{\ell} R_X K v d F = (\lambda n. (ldeg\text{-}p \text{ } R \text{ } (Vr \text{ } K \text{ } v) \text{ } X \text{ } d \text{ } (F \text{ } n)))]$$
$$Pseqh :: [(b, 'm) \text{ Ring-scheme}, b, -, b \Rightarrow ant, nat, nat \Rightarrow b] \Rightarrow nat \Rightarrow b$$

$$(\langle (6Pseqh \dots) \rangle [90,90,90,90,90,91]90) \text{ where}$$

$$Pseqh_R X K v d F = (\lambda n. (hdeg-p R (Vr K v) X (Suc d) (F n)))$$

```

lemma an-neq-minf1[simp]: $\forall n. \text{an } n \neq -\infty$ 
apply (rule allI) apply (simp add:an-def)
done

```

lemma (in Corps) *Pseqh-mem*: $\llbracket \text{valuation } K \ v; \text{PolynRg } R \ (Vr \ K \ v) \ X;$
 $F \ n \in \text{carrier } R; \forall n. \text{deg } R \ (Vr \ K \ v) \ X \ (F \ n) \leq an \ (Suc \ d) \rrbracket \implies$
 $(Pseqh_R \ X \ K \ v \ d \ F) \ n \in \text{carrier } R$
apply (frule *PolynRg.is-Ring*)
apply (frule *Vr-ring[of v]*)
apply (frule *PolynRg.subring[of R Vr K v X]*)
apply (frule *PolynRg.X-mem-R[of R Vr K v X]*)
apply (simp del:npow-suc add:Pseqh-def)
apply (rule *PolynRg.hdeg-p-mem*, assumption+, simp)
done

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$P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v^{(Vr \ K \ v) \ (an \ N)}) \ (ldeg\text{-}p \ R \ (Vr \ K \ v) \ X \ d \ p)$
apply (frule *PolynRg.is-Ring*)
apply (simp add:ldeg-p-def)
apply (frule *Vr-ring*[of *v*])
apply (frule *PolynRg.scf-d-pol*[of *R Vr K v X p Suc d*], assumption+,
(erule *conjE*)
apply (frule-tac *n = an N* in *vp-apow-ideal*[of *v*], simp)
apply (frule *PolynRg.P-mod-mod*[*THEN sym*, of *R Vr K v X vp K v (Vr K v) (an N)*
 $p \ scf\text{-}d \ R \ (Vr \ K \ v) \ X \ p \ (Suc \ d)$], assumption+, simp)
apply (subst *PolynRg.polyn-expr-short*[of *R Vr K v X*
 $scf\text{-}d \ R \ (Vr \ K \ v) \ X \ p \ (Suc \ d)$], assumption+, simp)
apply (subst *PolynRg.P-mod-mod*[*THEN sym*, of *R Vr K v X vp K v (Vr K v) (an N)*
 $polyn\text{-}expr \ R \ X \ d \ (d, \ snd \ (scf\text{-}d \ R \ (Vr \ K \ v) \ X \ p \ (Suc \ d)))$
 $(d, \ snd \ (scf\text{-}d \ R \ (Vr \ K \ v) \ X \ p \ (Suc \ d)))$], assumption+)
apply (subst *PolynRg.polyn-expr-short*[*THEN sym*], simp+,
simp add:*PolynRg.polyn-mem*)
apply (subst *pol-coeff-def*, rule *allI*, rule *impI*,
simp add:*PolynRg.pol-coeff-mem*)
apply simp+
done

lemma (in *Corps*) *PCauchy-hTr*: $\llbracket valuation \ K \ v; \ PolynRg \ R \ (Vr \ K \ v) \ X;$
 $p \in carrier \ R; \ deg \ R \ (Vr \ K \ v) \ X \ p \leq an \ (Suc \ d);$
 $P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v^{(Vr \ K \ v) \ (an \ N)}) \ p \rrbracket$
 $\implies P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v^{(Vr \ K \ v) \ (an \ N)}) \ (hdeg\text{-}p \ R \ (Vr \ K \ v) \ X \ (Suc$
 $d) \ p)$
apply (frule *PolynRg.is-Ring*)
apply (cut-tac *Vr-ring*[of *v*])
apply (frule *PolynRg.scf-d-pol*[of *R Vr K v X p Suc d*], assumption+)
apply (frule-tac *n = an N* in *vp-apow-ideal*[of *v*], simp)
apply (frule *PolynRg.P-mod-mod*[*THEN sym*, of *R Vr K v X*
 $vp \ K \ v \ (Vr \ K \ v) \ (an \ N) \ p \ scf\text{-}d \ R \ (Vr \ K \ v) \ X \ p \ (Suc \ d)$], assumption+,
simp+)
apply (subst *hdeg-p-def*)
apply (subst *PolynRg.monomial-P-mod-mod*[*THEN sym*, of *R Vr K v X*
 $vp \ K \ v \ (Vr \ K \ v) \ (an \ N) \ snd \ (scf\text{-}d \ R \ (Vr \ K \ v) \ X \ p \ (Suc \ d)) \ (Suc \ d)$
 $(snd \ (scf\text{-}d \ R \ (Vr \ K \ v) \ X \ p \ (Suc \ d)) \ (Suc \ d)) \cdot_{r \ R} \ (X^{\sim R} \ (Suc \ d))$
 $Suc \ d]$,
assumption+)
apply (rule *PolynRg.pol-coeff-mem*[of *R Vr K v X*
 $scf\text{-}d \ R \ (Vr \ K \ v) \ X \ p \ (Suc \ d) \ Suc \ d]$, assumption+, simp+)
done

lemma (in *Corps*) *v-ldeg-p-pOp*: $\llbracket valuation \ K \ v; \ PolynRg \ R \ (Vr \ K \ v) \ X;$
 $p \in carrier \ R; \ q \in carrier \ R; \ deg \ R \ (Vr \ K \ v) \ X \ p \leq an \ (Suc \ d);$
 $deg \ R \ (Vr \ K \ v) \ X \ q \leq an \ (Suc \ d) \rrbracket \implies$
 $(ldeg\text{-}p \ R \ (Vr \ K \ v) \ X \ d \ p) \pm_R (ldeg\text{-}p \ R \ (Vr \ K \ v) \ X \ d \ q) =$
 $ldeg\text{-}p \ R \ (Vr \ K \ v) \ X \ d \ (p \pm_R \ q)$

by (*simp add:PolynRg.ldeg-p-pOp*[of $R \text{ Vr } K \text{ v } X \text{ p } q \text{ d}$])

lemma (*in Corps*) *v-hdeg-p-pOp*: $\llbracket \text{valuation } K \text{ v}; \text{PolynRg } R \text{ (Vr } K \text{ v)} \text{ } X;$
 $p \in \text{carrier } R; q \in \text{carrier } R; \text{deg } R \text{ (Vr } K \text{ v)} \text{ } X \text{ p} \leq \text{an } (\text{Suc } d);$
 $\text{deg } R \text{ (Vr } K \text{ v)} \text{ } X \text{ q} \leq \text{an } (\text{Suc } d) \rrbracket \implies (\text{hdeg-p } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (Suc } d) \text{ p}) \pm_R$
 $(\text{hdeg-p } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (Suc } d) \text{ q}) = \text{hdeg-p } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (Suc } d) \text{ (p} \pm_R \text{ q)}$
by (*simp add:PolynRg.hdeg-p-pOp*[of $R \text{ Vr } K \text{ v } X \text{ p } q \text{ d}$])

lemma (*in Corps*) *v-ldeg-p-mOp*: $\llbracket \text{valuation } K \text{ v}; \text{PolynRg } R \text{ (Vr } K \text{ v)} \text{ } X;$
 $p \in \text{carrier } R; \text{deg } R \text{ (Vr } K \text{ v)} \text{ } X \text{ p} \leq \text{an } (\text{Suc } d) \rrbracket \implies$
 $-_a R (\text{ldeg-p } R \text{ (Vr } K \text{ v)} \text{ } X \text{ d } p) = \text{ldeg-p } R \text{ (Vr } K \text{ v)} \text{ } X \text{ d } (-_a R \text{ p})$
by (*simp add:PolynRg.ldeg-p-mOp*)

lemma (*in Corps*) *v-hdeg-p-mOp*: $\llbracket \text{valuation } K \text{ v}; \text{PolynRg } R \text{ (Vr } K \text{ v)} \text{ } X;$
 $p \in \text{carrier } R; \text{deg } R \text{ (Vr } K \text{ v)} \text{ } X \text{ p} \leq \text{an } (\text{Suc } d) \rrbracket \implies$
 $-_a R (\text{hdeg-p } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (Suc } d) \text{ p}) = \text{hdeg-p } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (Suc } d) \text{ } (-_a R \text{ p})$
by (*simp add:PolynRg.hdeg-p-mOp*)

lemma (*in Corps*) *PCauchy-lPCauchy*: $\llbracket \text{valuation } K \text{ v}; \text{PolynRg } R \text{ (Vr } K \text{ v)} \text{ } X;$
 $\forall n. F \text{ n} \in \text{carrier } R; \forall n. \text{deg } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (F } n) \leq \text{an } (\text{Suc } d);$
 $P\text{-mod } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (vp } K \text{ v}^{(\text{Vr } K \text{ v})} (\text{an } N)) \text{ (F } n \pm_R -_a R \text{ (F } m)) \rrbracket$
 $\implies P\text{-mod } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (vp } K \text{ v}^{(\text{Vr } K \text{ v})} (\text{an } N))$
 $((P\text{seq}_R \text{ X } K \text{ v } d \text{ F}) \text{ n}) \pm_R -_a R ((P\text{seq}_R \text{ X } K \text{ v } d \text{ F}) \text{ m}))$

apply (*frule PolynRg.is-Ring*)

apply (*cut-tac Vr-ring*[of v],

frule Ring.ring-is-ag[of R],

frule PolynRg.subring[of $R \text{ Vr } K \text{ v } X$])

apply (*simp add:Pseq-def*)

apply (*subst v-ldeg-p-mOp*[of $v \text{ R } X \text{ - } d$], *assumption+*,

simp, *simp*)

apply (*subst v-ldeg-p-pOp*[of $v \text{ R } X \text{ F } n \text{ } -_a R \text{ (F } m)$], *assumption+*,

simp, *rule aGroup.ag-mOp-closed*, *assumption*, *simp*, *simp*,

simp add:PolynRg.deg-minus-eq1)

apply (*rule PCauchy-lTr*[of $v \text{ R } X \text{ F } n \pm_R -_a R \text{ (F } m) \text{ d } N$],

assumption+,

rule aGroup.ag-pOp-closed[of $R \text{ F } n \text{ } -_a R \text{ (F } m)$], *assumption+*,

simp, *rule aGroup.ag-mOp-closed*, *assumption+*, *simp*)

apply (*frule PolynRg.deg-minus-eq1* [of $R \text{ Vr } K \text{ v } X \text{ F } m$], *simp*)

apply (*rule PolynRg.polyn-deg-add4*[of $R \text{ Vr } K \text{ v } X \text{ F } n$

$-_a R \text{ (F } m) \text{ Suc } d$], *assumption+*, *simp*,

rule aGroup.ag-mOp-closed, *assumption*, *simp+*)

done

lemma (*in Corps*) *PCauchy-hPCauchy*: $\llbracket \text{valuation } K \text{ v}; \text{PolynRg } R \text{ (Vr } K \text{ v)} \text{ } X;$
 $\forall n. F \text{ n} \in \text{carrier } R; \forall n. \text{deg } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (F } n) \leq \text{an } (\text{Suc } d);$
 $P\text{-mod } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (vp } K \text{ v}^{(\text{Vr } K \text{ v})} (\text{an } N)) \text{ (F } n \pm_R -_a R \text{ (F } m)) \rrbracket$
 $\implies P\text{-mod } R \text{ (Vr } K \text{ v)} \text{ } X \text{ (vp } K \text{ v}^{(\text{Vr } K \text{ v})} (\text{an } N))$
 $((P\text{seq}_R \text{ X } K \text{ v } d \text{ F}) \text{ n}) \pm_R -_a R ((P\text{seq}_R \text{ X } K \text{ v } d \text{ F}) \text{ m}))$


```

apply (frule PolynRg.is-Ring)
apply (frule Vr-ring[of v], frule Ring.ring-is-ag[of R],
      frule PolynRg.subring[of R Vr K v X],
      frule vp-apow-ideal[of v an N], simp)

apply (simp add:Pseqh-def,
      subst v-hdeg-p-mOp[of v R X F m d], assumption+,
      simp+)

apply (subst v-hdeg-p-pOp[of v R X F n -aR (F m)], assumption+,
      simp, rule aGroup.ag-mOp-closed, assumption, simp, simp,
      frule PolynRg.deg-minus-eq1 [of R Vr K v X F m],
      simp+)
apply (frule PCauchy-hTr[of v R X F n ±R -aR (F m) d N],
      assumption+,
      rule aGroup.ag-pOp-closed[of R F n -aR (F m)], assumption+,
      simp, rule aGroup.ag-mOp-closed, assumption+, simp)
apply (frule PolynRg.deg-minus-eq1 [of R Vr K v X F m], simp+)
apply (rule PolynRg.polyn-deg-add4[of R Vr K v X F n -aR (F m)
      Suc d], assumption+,
      simp, rule aGroup.ag-mOp-closed, assumption, simp+)
done

```

```

lemma (in Corps) Pseq-decompos:[valuation K v; PolynRg R (Vr K v) X;
  F n ∈ carrier R; deg R (Vr K v) X (F n) ≤ an (Suc d)]
  ⇒ F n = ((PseqlR X K v d F) n) ±R ((PseqhR X K v d F) n)
apply (frule PolynRg.is-Ring)
apply (simp del:npow-suc add:Pseql-def Pseqh-def)
apply (frule Vr-ring[of v])
apply (frule PolynRg.subring[of R Vr K v X])
apply (rule PolynRg.decompos-p[of R Vr K v X F n d], assumption+)
done

```

```

lemma (in Corps) deg-0-const:[valuation K v; PolynRg R (Vr K v) X;
  p ∈ carrier R; deg R (Vr K v) X p ≤ 0] ⇒ p ∈ carrier (Vr K v)
apply (frule Vr-ring[of v])
apply (frule PolynRg.subring)
apply (frule PolynRg.is-Ring)
apply (case-tac p = 0R, simp,
      simp add:Ring.Subring-zero-ring-zero[THEN sym],
      simp add:Ring.ring-zero)
apply (subst PolynRg.pol-of-deg0[THEN sym, of R Vr K v X p],
      assumption+)
apply (simp add:PolynRg.deg-an, simp only:an-0[THEN sym])
apply (simp only:ale-nat-le[of deg-n R (Vr K v) X p 0])
done

```

lemma (in Corps) monomial-P-limit: $\llbracket \text{valuation } K \ v; \text{Complete}_v \ K;$
 $\text{PolynRg } R \ (Vr \ K \ v) \ X; \forall n. f \ n \in \text{carrier } (Vr \ K \ v);$
 $\forall n. F \ n = (f \ n) \cdot_{\tau R} (X^{\sim R} \ d); \forall N. \exists M. \forall n \ m. M < n \wedge M < m \longrightarrow$
 $P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v) (Vr \ K \ v) \ (an \ N)) \ (F \ n \pm_R -_a R \ (F \ m)) \rrbracket \Longrightarrow$
 $\exists b \in \text{carrier } (Vr \ K \ v). \text{Plimit } R \ X \ K \ v \ F \ (b \cdot_{\tau R} (X^{\sim R} \ d))$
apply (frule PolynRg.is-Ring)
apply (frule Vr-ring[of v])
apply (frule PolynRg.subring[of R Vr K v X])
apply simp

apply (subgoal-tac Cauchy K v f)
apply (simp add:v-complete-def)
apply (drule-tac a = f in forall-spec)
apply (thin-tac $\forall N. \exists M. \forall n \ m. M < n \wedge M < m \longrightarrow$
 $P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v) (Vr \ K \ v) \ (an \ N))$
 $((f \ n) \cdot_{\tau R} (X^{\sim R} \ d) \pm_R -_a R \ (f \ m) \cdot_{\tau R} (X^{\sim R} \ d)), \text{assumption})$
apply (erule exE, erule conjE)
apply (subgoal-tac $b \in \text{carrier } (Vr \ K \ v)$)
apply (subgoal-tac $\text{Plimit } R \ X \ K \ v \ F \ (b \cdot_{\tau R} (X^{\sim R} \ d)), \text{blast}$)

apply (simp add:pol-limit-def)
apply (rule conjI)
apply (rule allI)
apply (rule Ring.ring-tOp-closed[of R], assumption)
apply (frule PolynRg.subring[of R Vr K v X])
apply (rule Ring.mem-subring-mem-ring[of R Vr K v], assumption+)
apply simp

apply (frule PolynRg.X-mem-R[of R Vr K v X])
apply (simp add:Ring.npClose)
apply (thin-tac $\forall n. F \ n = f \ n \cdot_{\tau R} X^{\sim R} \ d$)
apply (simp add:limit-def)
apply (rule allI)

apply (rotate-tac -2, drule-tac $x = N$ in spec)
apply (erule exE)
apply (subgoal-tac $\forall n > M. P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v) (Vr \ K \ v) \ (an \ N))$
 $((f \ n) \cdot_{\tau R} (X^{\sim R} \ d) \pm_R -_a R \ (b \cdot_{\tau R} (X^{\sim R} \ d))), \text{blast}$
apply (rule allI, rule impI)
apply (rotate-tac -2, drule-tac $x = n$ in spec, simp)
apply (drule-tac $x = n$ in spec)

apply (frule-tac $x = f \ n$ in Ring.mem-subring-mem-ring[of R Vr K v],
assumption+,
frule-tac $x = b$ in Ring.mem-subring-mem-ring[of R Vr K v],
assumption+)
apply (frule PolynRg.X-mem-R[of R Vr K v X])

```

apply (frule Ring.npClose[of R X d], assumption+)
apply (simp add:Ring.ring-inv1-1)
apply (frule Ring.ring-is-ag[of R],
  frule-tac  $x = b$  in aGroup.ag-mOp-closed[of R], assumption+)
apply (subst Ring.ring-distrib2[THEN sym, of R  $X^{\sim R} d$ ], assumption+)

apply (frule-tac  $n = an\ N$  in vp-apow-ideal[of v], simp)
apply (frule-tac  $I = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$  and  $c = f\ n\ \pm_R\ -_aR\ b$  and
   $p = (f\ n\ \pm_R\ -_aR\ b) \cdot_{rR}\ (X^{\sim R} d)$  in
  PolynRg.monomial-P-mod-mod[of R Vr K v X - - d], assumption+)
apply (simp add:Ring.Subring-minus-ring-minus[THEN sym])
apply (frule Ring.ring-is-ag[of Vr K v])
apply (frule-tac  $x = b$  in aGroup.ag-mOp-closed[of Vr K v], assumption+)
apply (simp only:Ring.Subring-pOp-ring-pOp[THEN sym])
apply (rule aGroup.ag-pOp-closed, assumption+) apply simp
apply (frule-tac  $I1 = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$  and  $c1 = f\ n\ \pm_R\ -_aR\ b$  and
   $p1 = (f\ n\ \pm_R\ -_aR\ b) \cdot_{rR}\ (X^{\sim R} d)$  in PolynRg.monomial-P-mod-mod[THEN
sym,
  of R Vr K v X - - d], assumption+)
apply (frule Ring.ring-is-ag[of Vr K v])
apply (frule-tac  $x = b$  in aGroup.ag-mOp-closed[of Vr K v], assumption+)
apply (simp only:Ring.Subring-minus-ring-minus[THEN sym])
apply (simp only:Ring.Subring-pOp-ring-pOp[THEN sym])
apply (rule aGroup.ag-pOp-closed, assumption+) apply simp apply simp
apply (simp only:Vr-mOp-f-mOp[THEN sym])
apply (frule Ring.ring-is-ag[of Vr K v])
apply (frule-tac  $x = b$  in aGroup.ag-mOp-closed[of Vr K v], assumption+)
apply (simp add:Vr-pOp-f-pOp[THEN sym])
apply (simp add:Ring.Subring-pOp-ring-pOp)
apply (simp add:Ring.Subring-minus-ring-minus)

apply (case-tac  $b = \mathbf{0}_K$ , simp add:Vr-0-f-0[THEN sym],
  simp add:Ring.ring-zero)
apply (frule-tac  $b = b$  in limit-val[of - f v], assumption+,
  rule allI,
  frule-tac  $x = j$  in spec, simp add:Vr-mem-f-mem,
  assumption+, erule exE)
apply (thin-tac  $\forall n. F\ n = f\ n \cdot_{rR}\ X^{\sim R} d$ ,
  thin-tac  $\forall N. \exists M. \forall n\ m. M < n \wedge M < m \longrightarrow$ 
   $P\text{-mod}\ R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))$ 
   $(f\ n \cdot_{rR}\ X^{\sim R} d \pm_R\ -_aR\ f\ m \cdot_{rR}\ X^{\sim R} d))$ 
apply (rotate-tac -1, drule-tac  $x = Suc\ N$  in spec, simp)
apply (drule-tac  $x = Suc\ N$  in spec)
apply (frule-tac  $x1 = f\ (Suc\ N)$  in val-pos-mem-Vr[THEN sym, of v],
  simp add:Vr-mem-f-mem, simp, simp add:val-pos-mem-Vr[of v])

apply (simp add:Cauchy-seq-def)
apply (simp add:Vr-mem-f-mem)

```

apply (*rule allI*)
apply (*rotate-tac* -3 , *frule-tac* $x = N$ **in** *spec*)

apply (*thin-tac* $\forall n. F\ n = f\ n \cdot_{rR} X^{\sim R} d$)

apply (*frule-tac* $n = an\ N$ **in** *vp-apow-ideal*[*of v*], *simp*)
apply (*drule-tac* $x = N$ **in** *spec*, *erule exE*)
apply (*subgoal-tac* $\forall n\ m. M < n \wedge M < m \longrightarrow$
 $f\ n \pm -_a (f\ m) \in vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$, *blast*)
apply (*rule allI*) +
apply (*rule impI*, *erule conjE*)
apply (*frule-tac* $I = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$ **and** $c = f\ n \pm -_a (f\ m)$ **and**
 $p = (f\ n \pm -_a (f\ m)) \cdot_{rR} (X^{\sim R} d)$ **in**
PolynRg.monomial-P-mod-mod[*of R Vr K v X - - - d*], *assumption+*)

apply (*frule-tac* $x = n$ **in** *spec*,
drule-tac $x = m$ **in** *spec*)
apply (*frule Ring.ring-is-ag*[*of Vr K v*],
simp add: Vr-mOp-f-mOp[*THEN sym*],
frule-tac $x = f\ m$ **in** *aGroup.ag-mOp-closed*[*of Vr K v*], *assumption+*,
simp add: Vr-pOp-f-pOp[*THEN sym*])
apply (*rule aGroup.ag-pOp-closed*, *assumption+*, *simp*)
apply *simp*
apply (*thin-tac* $(f\ n \pm -_a f\ m \in vp\ K\ v\ (Vr\ K\ v)\ (an\ N)) =$
 $P\text{-mod}\ R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))\ ((f\ n \pm -_a f\ m) \cdot_{rR} X^{\sim R} d)$)
apply (*rotate-tac* -3 , *drule-tac* $x = n$ **in** *spec*,
rotate-tac -1 , *drule-tac* $x = m$ **in** *spec*, *simp*)
apply (*frule-tac* $x = n$ **in** *spec*,
drule-tac $x = m$ **in** *spec*)
apply (*frule-tac* $x = f\ n$ **in** *Ring.mem-subring-mem-ring*[*of R Vr K v*],
assumption+,
frule-tac $x = f\ m$ **in** *Ring.mem-subring-mem-ring*[*of R Vr K v*],
assumption+,
frule Ring.ring-is-ag[*of R*],
frule-tac $x = f\ m$ **in** *aGroup.ag-mOp-closed*[*of R*], *assumption+*,
frule PolynRg.X-mem-R[*of R Vr K v X*],
frule Ring.npClose[*of R X d*], *assumption*)
apply (*simp add: Ring.ring-inv1-1*[*of R*],
frule-tac $y1 = f\ n$ **and** $z1 = -_aR\ f\ m$ **in** *Ring.ring-distrib2*[
THEN sym, *of R X^{\sim R} d*], *assumption+*, *simp*,
thin-tac $f\ n \cdot_{rR} X^{\sim R} d \pm_R (-_aR\ f\ m) \cdot_{rR} X^{\sim R} d =$
 $(f\ n \pm_R -_aR\ f\ m) \cdot_{rR} X^{\sim R} d$)
apply (*simp only: Ring.Subring-minus-ring-minus*[*THEN sym*, *of R Vr K v*])
apply (*frule Ring.subring-Ring*[*of R Vr K v*], *assumption*,
frule Ring.ring-is-ag[*of Vr K v*],
frule-tac $x = f\ m$ **in** *aGroup.ag-mOp-closed*[*of Vr K v*], *assumption+*)
apply (*simp add: Ring.Subring-pOp-ring-pOp*[*THEN sym*, *of R Vr K v*],
simp add: Vr-pOp-f-pOp, *simp add: Vr-mOp-f-mOp*)

done

lemma (in Corps) mPlimit-uniqueTr: $\llbracket \text{valuation } K \ v;$
 $\text{PolynRg } R \ (Vr \ K \ v) \ X; \forall n. f \ n \in \text{carrier } (Vr \ K \ v);$
 $\forall n. F \ n = (f \ n) \cdot_{rR} (X^{\sim R} \ d); \ c \in \text{carrier } (Vr \ K \ v);$
 $\text{Plimit } R \ X \ K \ v \ F \ (c \cdot_{rR} (X^{\sim R} \ d)) \rrbracket \implies \lim_{K \ v} f \ c$
apply (frule PolynRg.is-Ring,
simp add:pol-limit-def limit-def,
rule allI,
erule conjE,
rotate-tac -1, drule-tac $x = N$ in spec,
erule exE)
apply (subgoal-tac $\forall n. M < n \longrightarrow f \ n \pm -_a \ c \in \text{vp } K \ v \ (Vr \ K \ v) \ (an \ N), \text{blast}$)
apply (rule allI, rule impI,
rotate-tac -2, drule-tac $x = n$ in spec, simp,
drule-tac $x = n$ in spec,
drule-tac $x = n$ in spec,
thin-tac $\forall n. f \ n \cdot_{rR} X^{\sim R} \ d \in \text{carrier } R$)
apply (frule Vr-ring[of v],
frule PolynRg.X-mem-R[of R Vr K v X],
frule Ring.npClose[of R X d], assumption+,
frule PolynRg.subring[of R Vr K v X])
apply (frule-tac $x = c$ in Ring.mem-subring-mem-ring[of R Vr K v],
assumption+,
frule-tac $x = f \ n$ in Ring.mem-subring-mem-ring[of R Vr K v],
assumption+)
apply (simp add:Ring.ring-inv1-1,
frule Ring.ring-is-ag[of R],
frule aGroup.ag-mOp-closed[of R c], assumption+)
apply (simp add:Ring.ring-distrib2[THEN sym, of R $X^{\sim R} \ d - -_a \ c$],
simp add:Ring.Subring-minus-ring-minus[THEN sym],
frule Ring.ring-is-ag[of Vr K v],
frule aGroup.ag-mOp-closed[of Vr K v c], assumption+)
apply (simp add:Ring.Subring-pOp-ring-pOp[THEN sym],
frule-tac $x = f \ n$ in aGroup.ag-pOp-closed[of Vr K v -
 $-_a \ (Vr \ K \ v) \ c$], assumption+,
frule-tac $n = an \ N$ in vp-apow-ideal[of v], simp,
frule-tac $I1 = \text{vp } K \ v \ (Vr \ K \ v) \ (an \ N)$ and
 $cI = f \ n \pm_{(Vr \ K \ v)} -_a \ (Vr \ K \ v) \ c$ and $pI = (f \ n \pm_{(Vr \ K \ v)} -_a \ (Vr \ K \ v) \ c)$
 $\cdot_{rR} (X^{\sim R} \ d)$ in PolynRg.monomial-P-mod-mod[THEN sym, of R Vr K v
 $X - - - d$], assumption+, simp, simp)
apply (simp add:Vr-pOp-f-pOp, simp add:Vr-mOp-f-mOp)
done

lemma (in Corps) mono-P-limt-unique: $\llbracket \text{valuation } K \ v;$
 $\text{PolynRg } R \ (Vr \ K \ v) \ X; \forall n. f \ n \in \text{carrier } (Vr \ K \ v);$
 $\forall n. F \ n = (f \ n) \cdot_{rR} (X^{\sim R} \ d); \ b \in \text{carrier } (Vr \ K \ v); \ c \in \text{carrier } (Vr \ K \ v);$
 $\text{Plimit } R \ X \ K \ v \ F \ (b \cdot_{rR} (X^{\sim R} \ d)); \ \text{Plimit } R \ X \ K \ v \ F \ (c \cdot_{rR} (X^{\sim R} \ d)) \rrbracket \implies$

$$b \cdot_{rR} (X \sim^R d) = c \cdot_{rR} (X \sim^R d)$$
apply (*frule PolynRg.is-Ring*)
apply (*frule-tac mPlimit-uniqueTr*[of $v \ R \ X \ f \ F \ d \ b$], *assumption+*,
frule-tac mPlimit-uniqueTr[of $v \ R \ X \ f \ F \ d \ c$], *assumption+*)
apply (*frule Vr-ring*[of v],
frule PolynRg.subring[of $R \ Vr \ K \ v \ X$],
frule Vr-mem-f-mem[of $v \ b$], *assumption+*,
frule Vr-mem-f-mem[of $v \ c$], *assumption+*,
frule limit-unique[of $b \ f \ v \ c$])
apply (*rule allI*, *simp add: Vr-mem-f-mem*, *assumption+*, *simp*)
done

lemma (**in** *Corps*) *Plimit-deg*: $\llbracket$ *valuation* $K \ v$; *PolynRg* $R \ (Vr \ K \ v) \ X$;
 $\forall n. F \ n \in \text{carrier } R$; $\forall n. \text{deg } R \ (Vr \ K \ v) \ X \ (F \ n) \leq (an \ d)$;
 $p \in \text{carrier } R$; *Plimit* $R \ X \ K \ v \ F \ p \rrbracket \implies \text{deg } R \ (Vr \ K \ v) \ X \ p \leq (an \ d)$
apply (*frule PolynRg.is-Ring*, *frule Vr-ring*[of v])
apply (*case-tac* $p = \mathbf{0}_R$, *simp add:deg-def*)
apply (*rule contrapos-pp*, *simp+*,
simp add:aneg-le,
frule PolynRg.s-cf-expr[of $R \ Vr \ K \ v \ X \ p$], *assumption+*, (*erule conjE*)+,
frule PolynRg.s-cf-deg[of $R \ Vr \ K \ v \ X \ p$], *assumption+*,
frule PolynRg.pol-coeff-mem[of $R \ Vr \ K \ v \ X \ s\text{-cf } R \ (Vr \ K \ v) \ X \ p$
fst ($s\text{-cf } R \ (Vr \ K \ v) \ X \ p$)], *assumption+*, *simp*,
frule Vr-mem-f-mem[of $v \ \text{snd} \ (s\text{-cf } R \ (Vr \ K \ v) \ X \ p)$
fst ($s\text{-cf } R \ (Vr \ K \ v) \ X \ p$)], *assumption+*)

apply (*frule val-nonzero-noninf*[of v
 $\text{snd} \ (s\text{-cf } R \ (Vr \ K \ v) \ X \ p) \ (\text{fst} \ (s\text{-cf } R \ (Vr \ K \ v) \ X \ p))$], *assumption*,
simp add:Vr-0-f-0,
frule val-pos-mem-Vr[*THEN sym*, of $v \ \text{snd} \ (s\text{-cf } R \ (Vr \ K \ v) \ X \ p)$
fst ($s\text{-cf } R \ (Vr \ K \ v) \ X \ p$)], *assumption+*, *simp*,
frule value-in-aug-inf[of $v \ \text{snd} \ (s\text{-cf } R \ (Vr \ K \ v) \ X \ p)$
fst ($s\text{-cf } R \ (Vr \ K \ v) \ X \ p$)], *assumption+*,
cut-tac mem-ant[of $v \ (\text{snd} \ (s\text{-cf } R \ (Vr \ K \ v) \ X \ p)$
fst ($s\text{-cf } R \ (Vr \ K \ v) \ X \ p$))], *simp add:aug-inf-def*,
erule exE, *simp*, *simp only:ant-0[THEN sym]*, *simp only:ale-zle*,
frule-tac $z = z$ **in** *zpos-nat*, *erule exE*, *simp*,
thin-tac $z = \text{int } n$)

apply (*simp add:pol-limit-def*)
apply (*rotate-tac* 5, *drule-tac* $x = \text{Suc } n$ **in** *spec*)
apply (*erule exE*)
apply (*rotate-tac* -1,
drule-tac $x = \text{Suc } M$ **in** *spec*, *simp del:npow-suc*,
drule-tac $x = \text{Suc } M$ **in** *spec*,
drule-tac $x = \text{Suc } M$ **in** *spec*)

apply (*frule* *PolynRg.polyn-minus*[*of* *R Vr K v X s-cf R (Vr K v) X p*
fst (s-cf R (Vr K v) X p)], *assumption+*, *simp*,
frule *PolynRg.minus-pol-coeff*[*of* *R Vr K v X s-cf R (Vr K v) X p*],
assumption+, *drule sym*,
frule-tac *x = deg R (Vr K v) X (F (Suc M))* **in** *ale-less-trans*[*of* -
an d deg R (Vr K v) X p], *assumption+*,
frule-tac *p = F (Suc M)* **and** *d = deg-n R (Vr K v) X p* **in**
PolynRg.pol-expr-edeg[*of* *R Vr K v X*], *assumption+*,
frule-tac *x = deg R (Vr K v) X (F (Suc M))* **and**
y = deg R (Vr K v) X p **in** *aless-imp-le*,
subst *PolynRg.deg-an*[*THEN sym*, *of* *R Vr K v X p*], *assumption+*,
erule exE, (*erule conjE*)+,
frule-tac *c = f* **in** *PolynRg.polyn-add1*[*of* *R Vr K v X -*
(fst (s-cf R (Vr K v) X p), $\lambda j. -_a Vr K v \text{ snd } (s-cf R (Vr K v) X p) j)$)],
assumption+, *simp*,
thin-tac $-_a R p = \text{polyn-expr } R X (\text{fst } (s-cf R (Vr K v) X p))$
 $(\text{fst } (s-cf R (Vr K v) X p), \lambda j. -_a Vr K v \text{ snd } (s-cf R (Vr K v) X p) j)$,
thin-tac *polyn-expr R X (fst (s-cf R (Vr K v) X p))*
 $(s-cf R (Vr K v) X p) = p$,
thin-tac *F (Suc M) = polyn-expr R X (fst (s-cf R (Vr K v) X p)) f*,
thin-tac *polyn-expr R X (fst (s-cf R (Vr K v) X p)) f $\pm_R$*
polyn-expr R X (fst (s-cf R (Vr K v) X p))
 $(\text{fst } (s-cf R (Vr K v) X p), \lambda j. -_a Vr K v \text{ snd } (s-cf R (Vr K v) X p) j) =$
 $\text{polyn-expr } R X (\text{fst } (s-cf R (Vr K v) X p)) (\text{add-cf } (Vr K v) f$
 $(\text{fst } (s-cf R (Vr K v) X p), \lambda j. -_a Vr K v \text{ snd } (s-cf R (Vr K v) X p) j)))$

apply (*frule-tac* *n = an (Suc n)* **in** *vp-apow-ideal*[*of* *v*], *simp*,
frule-tac *p1 = polyn-expr R X (fst (s-cf R (Vr K v) X p)) (add-cf (Vr K v) f*
 $(\text{fst } (s-cf R (Vr K v) X p), \lambda j. -_a Vr K v \text{ snd } (s-cf R (Vr K v) X p) j))$ **and**
 $I1 = vp K v (Vr K v) (an (Suc n))$ **and** *c1 = add-cf (Vr K v) f (fst*
 $(s-cf R (Vr K v) X p), \lambda j. -_a Vr K v \text{ snd } (s-cf R (Vr K v) X p) j)$ **in**
PolynRg.P-mod-mod[*THEN sym*, *of* *R Vr K v X*], *assumption+*,
rule *PolynRg.polyn-mem*[*of* *R Vr K v X*], *assumption+*,
rule *PolynRg.add-cf-pol-coeff*[*of* *R Vr K v X*], *assumption+*,
simp *add:PolynRg.add-cf-len*,
rule *PolynRg.add-cf-pol-coeff*[*of* *R Vr K v X*], *assumption+*,
simp *add:PolynRg.add-cf-len*,
simp *add:PolynRg.add-cf-len*)

apply (*drule-tac* *x = fst (s-cf R (Vr K v) X p)* **in** *spec*, *simp*,
thin-tac *P-mod R (Vr K v) X (vp K v (Vr K v) (an (Suc n)))*
 $(\text{polyn-expr } R X (\text{fst } (s-cf R (Vr K v) X p)) (\text{add-cf } (Vr K v) f$
 $(\text{fst } (s-cf R (Vr K v) X p), \lambda j. -_a Vr K v \text{ snd } (s-cf R (Vr K v) X p) j)))$,
simp *add:add-cf-def*)

apply (*frule-tac* *p = polyn-expr R X (fst (s-cf R (Vr K v) X p)) f* **and**
c = f **and** *j = fst f* **in** *PolynRg.pol-len-gt-deg*[*of* *R Vr K v X*],
assumption+, *simp*, *drule sym*, *simp* *add:PolynRg.deg-an*) **apply** *simp*
apply (*rotate-tac* -4 , *drule sym*, *simp*)

```

apply (frule Ring.ring-is-ag[of Vr K v],
  frule-tac  $x = \text{snd } (s\text{-cf } R \text{ (Vr K v) } X p) \text{ (fst } f) \text{ in}$ 
  aGroup.ag-mOp-closed[of Vr K v], assumption+,
  simp add:aGroup.ag-l-zero)
apply (frule-tac  $I = v p K v \text{ (Vr K v) (an (Suc n)) and}$ 
   $x = -_a V_r K v \text{ snd } (s\text{-cf } R \text{ (Vr K v) } X p) \text{ (fst } f) \text{ in}$ 
  Ring.ideal-inv1-closed[of Vr K v], assumption+)
apply (simp add:aGroup.ag-inv-inv)
apply (frule-tac  $n = \text{an (Suc n) and } x = \text{snd } (s\text{-cf } R \text{ (Vr K v) } X p) \text{ (fst } f) \text{ in}$ 
  in n-value-x-1[of v], simp+)
apply (frule-tac  $x = \text{snd } (s\text{-cf } R \text{ (Vr K v) } X p) \text{ (fst } f) \text{ in}$ 
  n-val-le-val[of v], assumption+, simp add:ant-int)
apply (drule-tac  $i = \text{an (Suc n) and}$ 
   $j = \text{n-val } K v \text{ (snd } (s\text{-cf } R \text{ (Vr K v) } X p) \text{ (fst } f)) \text{ and}$ 
   $k = v \text{ (snd } (s\text{-cf } R \text{ (Vr K v) } X p) \text{ (fst } f)) \text{ in ale-trans,}$ 
  assumption+)
apply (simp add:ant-int ale-natle)
done

lemma (in Corps) Plimit-deg1:  $\llbracket \text{valuation } K v; \text{Ring } R; \text{PolynRg } R \text{ (Vr K v) } X;$ 
 $\forall n. F n \in \text{carrier } R; \forall n. \text{deg } R \text{ (Vr K v) } X (F n) \leq ad;$ 
 $p \in \text{carrier } R; \text{Plimit } R X K v F p \rrbracket \implies \text{deg } R \text{ (Vr K v) } X p \leq ad$ 
apply (frule Vr-ring[of v])
apply (case-tac  $\forall n. F n = \mathbf{0}_R$ )
apply (frule Plimit-deg[of v R X F 0 p], assumption+,
  rule allI, simp add:deg-def, assumption+)
apply (case-tac  $p = \mathbf{0}_R$ , simp add:deg-def,
  frule PolynRg.nonzero-deg-pos[of R Vr K v X p], assumption+,
  simp,
  frule PolynRg.pols-const[of R Vr K v X p], assumption+,
  simp,
  frule PolynRg.pols-const[of R Vr K v X p], assumption+,
  simp add:ale-refl)
apply (subgoal-tac  $p = \mathbf{0}_R$ , simp)
apply (thin-tac  $p \neq \mathbf{0}_R$ )
apply (rule contrapos-pp, simp+)

apply (frule n-val-valuation[of v])
apply (frule val-nonzero-z[of n-val K v p])
apply (simp add:Vr-mem-f-mem)
apply (frule PolynRg.subring[of R Vr K v X])
apply (simp only:Ring.Subring-zero-ring-zero[THEN sym, of R Vr K v])
apply (simp add:Vr-0-f-0, erule exE)
apply (frule val-pos-mem-Vr[THEN sym, of v p])
apply (simp add:Vr-mem-f-mem, simp)
apply (frule val-pos-n-val-pos[of v p])
apply (simp add:Vr-mem-f-mem, simp)
apply (simp add:ant-0[THEN sym])
apply (frule-tac  $z = z \text{ in } \text{zpos-nat, erule exE}$ )

```


apply (*unfold pol-limit-def*, *erule conjE*)
apply (*rotate-tac* -1 , *drule-tac* $x = \text{Suc } n$ **in** *spec*)
apply (*subgoal-tac* $\neg (\exists M. \forall m. M < m \longrightarrow$
 $P\text{-mod } R (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ (\text{Suc } n)))\ (F\ m\ \pm_R\ -_aR\ p)))$
apply *blast*
apply (*thin-tac* $\exists M. \forall m. M < m \longrightarrow$
 $P\text{-mod } R (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ (\text{Suc } n)))\ (F\ m\ \pm_R\ -_aR\ p))$
apply *simp*
apply (*subgoal-tac* $M < (\text{Suc } M) \wedge$
 $\neg P\text{-mod } R (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ (\text{Suc } n)))\ (\mathbf{0}_R\ \pm_R\ -_aR\ p))$
apply *blast*
apply *simp*
apply (*frule* *Ring.ring-is-ag*[*of* R])
apply (*frule* *aGroup.ag-mOp-closed*[*of* $R\ p$], *assumption*)
apply (*simp* *add:aGroup.ag-l-zero*)
apply (*frule* *Ring.ring-is-ag*[*of* $Vr\ K\ v$])
apply (*frule* *aGroup.ag-mOp-closed*[*of* $Vr\ K\ v\ p$], *assumption*)
apply (*frule-tac* $n = an\ (\text{Suc } n)$ **in** *vp-apow-ideal*[*of* v], *simp*)
apply (*frule* *PolynRg.subring*[*of* $R\ Vr\ K\ v\ X$])
apply (*simp* *add:Ring.Subring-minus-ring-minus*[*THEN* *sym*, *of* $R\ Vr\ K\ v$])
apply (*simp* *add:PolynRg.P-mod-coeffTr*[*of* $R\ Vr\ K\ v\ X\ -\ -_a(Vr\ K\ v)\ p$])
apply (*rule* *contrapos-pp*, *simp*+)

apply (*frule-tac* $I = vp\ K\ v\ (Vr\ K\ v)\ (an\ (\text{Suc } n))$ **in**
 $Ring.ideal-inv1-closed$ [*of* $Vr\ K\ v\ -\ -_a(Vr\ K\ v)\ p$], *assumption*+)

apply (*simp* *add:aGroup.ag-inv-inv*)
apply (*frule-tac* $n = an\ (\text{Suc } n)$ **in** *n-value-x-I*[*of* $v\ -\ p$], *simp*)
apply *assumption*
apply *simp*
apply (*simp* *add:ant-int*, *simp* *add:ale-natle*)

apply (*fold pol-limit-def*)
apply (*case-tac* $ad = \infty$, *simp*)
apply *simp* **apply** (*erule exE*)
apply (*subgoal-tac* $0 \leq ad$)
apply (*frule* *Plimit-deg*[*of* $v\ R\ X\ F\ na\ ad\ p$], *assumption*+)

apply (*simp* *add:an-na*) +
apply (*drule-tac* $x = n$ **in** *spec*,
 $drule-tac\ x = n$ **in** *spec*)

apply (*frule-tac* $p = F\ n$ **in** *PolynRg.nonzero-deg-pos*[*of* $R\ Vr\ K\ v\ X$],
 $assumption$ +)

apply (*rule-tac* $j = \text{deg } R\ (Vr\ K\ v)\ X\ (F\ n)$ **in** *ale-trans*[*of* $0\ -\ ad$],
 $assumption$ +)

done

lemma (**in** *Corps*) *Plimit-ldeg*: $\llbracket valuation\ K\ v; PolynRg\ R\ (Vr\ K\ v)\ X;$
 $\forall n. F\ n \in carrier\ R; p \in carrier\ R;$

$\forall n. \deg R (Vr K v) X (F n) \leq an (Suc d);$
 $Plimit\ R\ X\ K\ v\ F\ p \Longrightarrow Plimit\ R\ X\ K\ v\ (Pseq\ R\ X\ K\ v\ d\ F)$
 $(ldeg-p\ R\ (Vr\ K\ v)\ X\ d\ p)$
apply (frule Vr-ring[of v], frule PolynRg.is-Ring,
frule Ring.ring-is-ag[of R])
apply (frule Plimit-deg[of v R X F Suc d p], assumption+)
apply (simp add:Pseq-def, simp add:pol-limit-def)
apply (rule conjI, rule allI)
apply (rule PolynRg.ldeg-p-mem, assumption+, simp+)
apply (rule allI)
apply (rotate-tac -5, drule-tac $x = N$ in spec, erule exE)
apply (subgoal-tac $\forall m > M. P\text{-mod } R (Vr K v) X (vp K v (Vr K v) (an N))$
 $(ldeg-p\ R\ (Vr K v) X d (F m) \pm_R -_a R (ldeg-p\ R\ (Vr K v) X d p)),$
blast)
apply (rule allI, rule impI)
apply (rotate-tac -2,
frule-tac $x = m$ in spec,
thin-tac $\forall m. M < m \longrightarrow$
 $P\text{-mod } R (Vr K v) X (vp K v (Vr K v) (an N)) (F m \pm_R -_a R p),$
simp)
apply (subst v-ldeg-p-mOp[of v R X - d], assumption+)
apply (subst v-ldeg-p-pOp[of v R X - $-_a R$ p], assumption+)
apply (simp, rule aGroup.ag-mOp-closed, assumption, simp, simp)
apply (frule PolynRg.deg-minus-eq1 [THEN sym, of R Vr K v X p],
assumption+)
apply simp
apply (rule-tac $p = F m \pm_R -_a R p$ and $N = N$ in PCauchy-lTr[of v
R X - d], assumption+)
apply (rule-tac $x = F m$ in aGroup.ag-pOp-closed[of R - $-_a R$ p],
assumption+)
apply (simp, rule aGroup.ag-mOp-closed, assumption+)
apply (frule PolynRg.deg-minus-eq1 [of R Vr K v X p], assumption+)
apply (rule PolynRg.polyn-deg-add4[of R Vr K v X - $-_a R$ p Suc d],
assumption+)
apply (simp, rule aGroup.ag-mOp-closed, assumption, simp+)
done

lemma (in Corps) Plimit-hdeg: $\llbracket valuation\ K\ v; PolynRg\ R\ (Vr\ K\ v)\ X;$
 $\forall n. F n \in carrier\ R; \forall n. \deg R (Vr K v) X (F n) \leq an (Suc d);$
 $p \in carrier\ R; Plimit\ R\ X\ K\ v\ F\ p \Longrightarrow$
 $Plimit\ R\ X\ K\ v\ (Pseqh\ R\ X\ K\ v\ d\ F) (hdeg-p\ R\ (Vr\ K\ v)\ X\ (Suc\ d)\ p)$
apply (frule Vr-ring[of v], frule PolynRg.is-Ring,
frule Ring.ring-is-ag[of R])
apply (frule Plimit-deg[of v R X F Suc d p], assumption+)
apply (simp add:Pseqh-def, simp add:pol-limit-def)
apply (rule conjI, rule allI)
apply (rule PolynRg.hdeg-p-mem, assumption+, simp+)
apply (rule allI)
apply (rotate-tac -5, drule-tac $x = N$ in spec)

```

apply (erule exE)
apply (subgoal-tac  $\forall m > M. P\text{-mod } R (Vr K v) X (vp K v (Vr K v) (an N))$ 
  ( $hdeg\text{-}p R (Vr K v) X (Suc d) (F m) \pm_R -_a R (hdeg\text{-}p R (Vr K v) X (Suc d)$ 
   $p)$ ),
  blast)
apply (rule allI, rule impI)
apply (rotate-tac -2,
  drule-tac  $x = m$  in spec, simp)
apply (subst v-hdeg-p-mOp[of  $v R X - d$ ], assumption+)
apply (subst v-hdeg-p-pOp[of  $v R X - -_a R p$ ], assumption+)
apply (simp, rule aGroup.ag-mOp-closed, assumption, simp, simp)
apply (frule PolynRg.deg-minus-eq1 [THEN sym, of  $R Vr K v X p$ ], assumption+)
apply simp
apply (rule-tac  $p = F m \pm_R -_a R p$  and  $N = N$  in PCauchy-hTr[of  $v R X - d$ ],
  assumption+)
apply (rule-tac  $x = F m$  in aGroup.ag-pOp-closed[of  $R - -_a R p$ ],
  assumption+)
apply (simp, rule aGroup.ag-mOp-closed, assumption+)
apply (frule PolynRg.deg-minus-eq1 [of  $R Vr K v X p$ ], assumption+)
apply (rule PolynRg.polyn-deg-add4[of  $R Vr K v X - -_a R p Suc d$ ],
  assumption+)
apply (simp, rule aGroup.ag-mOp-closed, assumption, simp+)
done

```

```

lemma (in Corps)  $P\text{-limit-uniqueTr} : \llbracket valuation K v; PolynRg R (Vr K v) X \rrbracket \implies$ 
 $\forall F. ((\forall n. F n \in carrier R) \wedge (\forall n. deg R (Vr K v) X (F n) \leq (an d)) \longrightarrow$ 
 $(\forall p1 p2. p1 \in carrier R \wedge p2 \in carrier R \wedge Plimit R X K v F p1 \wedge$ 
 $Plimit R X K v F p2 \longrightarrow p1 = p2))$ 
apply (frule PolynRg.is-Ring)
apply (induct-tac d)
apply (rule allI, rule impI, (rule allI)+, rule impI)
apply (erule conjE)+
apply (subgoal-tac  $\forall n. F n \in carrier (Vr K v)$ )
apply (frule Vr-ring[of  $v$ ])
apply (frule PolynRg.X-mem-R[of  $R Vr K v X$ ])
apply (frule-tac  $f = F$  and  $F = F$  and  $d = 0$  and  $b = p1$  and  $c = p2$  in
  mono-P-limt-unique[of  $v R X$ ], assumption+)
apply (rule allI, drule-tac  $x = n$  in spec,
  simp add: Ring.ring-r-one)
apply (frule-tac  $F = F$  and  $p = p1$  in Plimit-deg[of  $v R X - 0$ ],
  assumption+, simp add: deg-0-const,
  frule-tac  $F = F$  and  $p = p2$  in Plimit-deg[of  $v R X - 0$ ],
  assumption+, simp add: deg-0-const)
apply (simp add: Ring.ring-r-one)+
apply (simp add: deg-0-const)

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apply (rename-tac d)
apply (rule allI, rule impI)

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apply (erule conjE)
apply ((rule allI)+, rule impI, (erule conjE)+)
apply (frule-tac  $F = F$  and  $p = p1$  and  $d = d$  in Plimit-ldeg[of v R X],
      assumption+,
      frule-tac  $F = F$  and  $p = p2$  and  $d = d$  in Plimit-ldeg[of v R X],
      assumption+,
      frule-tac  $F = F$  and  $p = p1$  and  $d = d$  in Plimit-hdeg[of v R X],
      assumption+,
      frule-tac  $F = F$  and  $p = p2$  and  $d = d$  in Plimit-hdeg[of v R X],
      assumption+)
apply (frule-tac  $a = Pseq1\ R\ X\ K\ v\ d\ F$  in forall-spec)
apply (rule conjI)
apply (rule allI)
apply (rule Pseq1-mem, assumption+, simp)
apply (rule allI, simp)
apply (rule allI)
apply (subst Pseq1-def)
apply (rule-tac  $p = F\ n$  and  $d = d$  in PolynRg.deg-ldeg-p[of R Vr K v X],
      assumption+) apply (simp add: Vr-ring)

apply simp
apply (thin-tac  $\forall F. (\forall n. F\ n \in \text{carrier } R) \wedge$ 
       $(\forall n. \text{deg } R\ (Vr\ K\ v)\ X\ (F\ n) \leq an\ d) \longrightarrow$ 
       $(\forall p1\ p2.$ 
         $p1 \in \text{carrier } R \wedge$ 
         $p2 \in \text{carrier } R \wedge$ 
         $Plimit\ R\ X\ K\ v\ F\ p1 \wedge Plimit\ R\ X\ K\ v\ F\ p2 \longrightarrow$ 
         $p1 = p2))$ 
apply (frule Vr-ring[of v])
apply (frule-tac  $F = F$  and  $d = Suc\ d$  and  $p = p1$  in
      Plimit-deg[of v R X], assumption+,
      frule-tac  $F = F$  and  $d = Suc\ d$  and  $p = p2$  in
      Plimit-deg[of v R X], assumption+)
apply (subgoal-tac (ldeg-p  $R\ (Vr\ K\ v)\ X\ d\ p1$ ) = (ldeg-p  $R\ (Vr\ K\ v)\ X\ d\ p2$ ))
apply (subgoal-tac hdeg-p  $R\ (Vr\ K\ v)\ X\ (Suc\ d)\ p1$  =
      hdeg-p  $R\ (Vr\ K\ v)\ X\ (Suc\ d)\ p2$ )

apply (frule-tac  $p = p1$  and  $d = d$  in PolynRg.decompos-p[of R Vr K v X],
      assumption+,
      frule-tac  $p = p2$  and  $d = d$  in PolynRg.decompos-p[of R Vr K v X],
      assumption+)

apply simp
apply (thin-tac  $Plimit\ R\ X\ K\ v\ Pseq1\ R\ X\ K\ v\ d\ F\ (ldeg-p\ R\ (Vr\ K\ v)\ X\ d\ p1),$ 
      thin-tac  $Plimit\ R\ X\ K\ v\ Pseq1\ R\ X\ K\ v\ d\ F\ (ldeg-p\ R\ (Vr\ K\ v)\ X\ d\ p2),$ 
      thin-tac  $\forall p1\ p2. p1 \in \text{carrier } R \wedge p2 \in \text{carrier } R \wedge$ 
       $Plimit\ R\ X\ K\ v\ Pseq1\ R\ X\ K\ v\ d\ F\ p1 \wedge$ 
       $Plimit\ R\ X\ K\ v\ Pseq1\ R\ X\ K\ v\ d\ F\ p2 \longrightarrow p1 = p2)$ 
apply (simp only:hdeg-p-def)
apply (rule-tac  $f = \lambda j. \text{snd } (scf-d\ R\ (Vr\ K\ v)\ X\ (F\ j)\ (Suc\ d))\ (Suc\ d)$ 
      and  $F = Pseqh\ R\ X\ K\ v\ d\ F$ )

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and  $b = (\text{snd } (\text{scf-d } R \text{ (Vr } K \text{ v) } X \text{ p1 } (\text{Suc } d))) (\text{Suc } d)$ 
and  $d = \text{Suc } d$  and  $c = \text{snd } (\text{scf-d } R \text{ (Vr } K \text{ v) } X \text{ p2 } (\text{Suc } d)) (\text{Suc } d)$  in
   $\text{mono-}P\text{-limit-unique}[\text{of } v \text{ } R \text{ } X], \text{assumption+})$ 
apply (rule allI)
apply (frule-tac  $p = F \text{ } n$  and  $d = \text{Suc } d$  in
   $\text{PolynRg.scf-d-pol}[\text{of } R \text{ Vr } K \text{ v } X]$ )
apply (simp del:npow-suc) +
apply (frule-tac  $c = \text{scf-d } R \text{ (Vr } K \text{ v) } X (F \text{ } n) (\text{Suc } d)$  and
   $j = \text{Suc } d$  in  $\text{PolynRg.pol-coeff-mem}[\text{of } R \text{ Vr } K \text{ v } X]$ )
apply simp apply (simp del:npow-suc) +
apply (rule allI)
apply (frule-tac  $c = \text{scf-d } R \text{ (Vr } K \text{ v) } X (F \text{ } n) (\text{Suc } d)$  and
   $j = \text{Suc } d$  in  $\text{PolynRg.pol-coeff-mem}[\text{of } R \text{ Vr } K \text{ v } X]$ )
apply (frule-tac  $p = F \text{ } n$  and  $d = \text{Suc } d$  in  $\text{PolynRg.scf-d-pol}[\text{of } R$ 
   $\text{Vr } K \text{ v } X], (\text{simp del:npow-suc}) +$ )
apply (cut-tac  $p = F \text{ } n$  and  $d = \text{Suc } d$  in  $\text{PolynRg.scf-d-pol}[\text{of } R$ 
   $\text{Vr } K \text{ v } X], (\text{simp del:npow-suc}) +$ )

apply (subst Pseqh-def) apply (simp only:hdeg-p-def)
apply (frule-tac  $p = p1$  and  $d = \text{Suc } d$  in  $\text{PolynRg.scf-d-pol}[\text{of } R \text{ Vr } K \text{ v } X],$ 
  assumption+)
apply (rule-tac  $c = \text{scf-d } R \text{ (Vr } K \text{ v) } X \text{ p1 } (\text{Suc } d)$  and
   $j = \text{Suc } d$  in  $\text{PolynRg.pol-coeff-mem}[\text{of } R \text{ Vr } K \text{ v } X], \text{assumption+},$ 
  simp, simp)
apply (frule-tac  $p = p2$  and  $d = \text{Suc } d$  in  $\text{PolynRg.scf-d-pol}[\text{of } R \text{ Vr } K \text{ v } X],$ 
  assumption+)
apply (rule-tac  $c = \text{scf-d } R \text{ (Vr } K \text{ v) } X \text{ p2 } (\text{Suc } d)$  and
   $j = \text{Suc } d$  in  $\text{PolynRg.pol-coeff-mem}[\text{of } R \text{ Vr } K \text{ v } X], \text{assumption+},$ 
  simp, simp) apply simp apply simp
apply (rotate-tac  $-4,$ 
  drule-tac  $x = \text{ldeg-p } R \text{ (Vr } K \text{ v) } X \text{ d } p1$  in spec,
  rotate-tac  $-1,$ 
  drule-tac  $x = \text{ldeg-p } R \text{ (Vr } K \text{ v) } X \text{ d } p2$  in spec)
apply (simp add:PolynRg.ldeg-p-mem)
done

```

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lemma (in Corps)  $P\text{-limit-unique}[\llbracket \text{valuation } K \text{ } v; \text{Complete}_v \text{ } K;$ 
   $\text{PolynRg } R \text{ (Vr } K \text{ v) } X; \forall n. F \text{ } n \in \text{carrier } R;$ 
   $\forall n. \text{deg } R \text{ (Vr } K \text{ v) } X (F \text{ } n) \leq (\text{an } d); p1 \in \text{carrier } R; p2 \in \text{carrier } R;$ 
   $\text{Plimit } R \text{ } X \text{ } K \text{ } v \text{ } F \text{ } p1; \text{Plimit } R \text{ } X \text{ } K \text{ } v \text{ } F \text{ } p2 \rrbracket \implies p1 = p2$ 
apply (frule  $P\text{-limit-uniqueTr}[\text{of } v \text{ } R \text{ } X \text{ } d], \text{assumption+})$ 
apply blast
done

```

```

lemma (in Corps)  $P\text{-limitTr}[\llbracket \text{valuation } K \text{ } v; \text{Complete}_v \text{ } K; \text{PolynRg } R \text{ (Vr } K \text{ v) } X \rrbracket$ 
 $\implies \forall F. ((\forall n. F \text{ } n \in \text{carrier } R) \wedge (\forall n. \text{deg } R \text{ (Vr } K \text{ v) } X (F \text{ } n) \leq (\text{an } d)) \wedge$ 
   $(\forall N. \exists M. \forall n \text{ } m. M < n \wedge M < m \implies$ 
   $P\text{-mod } R \text{ (Vr } K \text{ v) } X (vp \text{ } K \text{ } v \text{ (Vr } K \text{ v) } (\text{an } N)) (F \text{ } n \pm_R \text{ } {}^{-a} R (F \text{ } m))) \implies$ 

```

$(\exists p \in \text{carrier } R. \text{Plimit } R \ X \ K \ v \ F \ p))$
apply (frule PolynRg.is-Ring)
apply (frule Vr-ring[of v])
apply (induct-tac d)

apply simp
apply (rule allI, rule impI, (erule conjE)+)
apply (frule-tac $F = F$ **and** $f = F$ **in** monomial-P-limit[of v R X - - 0],
 assumption+)
apply (rule allI)
apply (rotate-tac 5, drule-tac $x = n$ **in** spec)
apply (simp add:deg-0-const)
apply (rule allI)
apply (drule-tac $x = n$ **in** spec,
 thin-tac $\forall n. \text{deg } R \ (Vr \ K \ v) \ X \ (F \ n) \leq 0$)
apply (frule PolynRg.X-mem-R[of R Vr K v X],
 frule PolynRg.is-Ring,
 simp add:Ring.ring-r-one, assumption)

apply (erule bexE)
apply (frule PolynRg.subring[of R Vr K v X])
apply (cut-tac $x = b$ **in** Ring.mem-subring-mem-ring[of R Vr K v])
apply (frule PolynRg.X-mem-R[of R Vr K v X], assumption+)
apply (simp add:Ring.ring-r-one, blast)

apply (rule allI, rule impI) **apply** (rename-tac d F)
apply (erule conjE)+
apply (subgoal-tac $(\forall n. (Pseq1 \ R \ X \ K \ v \ d \ F) \ n \in \text{carrier } R) \wedge$
 $(\forall n. \text{deg } R \ (Vr \ K \ v) \ X \ ((Pseq1 \ R \ X \ K \ v \ d \ F) \ n) \leq an \ d) \wedge$
 $(\forall N. \exists M. \forall n \ m. M < n \wedge M < m \longrightarrow$
 $P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v \ (Vr \ K \ v) \ (an \ N))$
 $((Pseq1 \ R \ X \ K \ v \ d \ F) \ n \pm_R -_a R \ ((Pseq1 \ R \ X \ K \ v \ d \ F) \ m))))$)
apply (frule-tac $a = Pseq1 \ R \ X \ K \ v \ d \ F$ **in** forall-spec, assumption)
apply (erule bexE)
apply (thin-tac $\forall F. (\forall n. F \ n \in \text{carrier } R) \wedge$
 $(\forall n. \text{deg } R \ (Vr \ K \ v) \ X \ (F \ n) \leq an \ d) \wedge$
 $(\forall N. \exists M. \forall n \ m. M < n \wedge M < m \longrightarrow$
 $P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v \ (Vr \ K \ v) \ (an \ N)) \ (F \ n \pm_R -_a R \ (F \ m))) \longrightarrow$
 $(\exists p \in \text{carrier } R. \text{Plimit } R \ X \ K \ v \ F \ p),$
 thin-tac $(\forall n. (Pseq1 \ R \ X \ K \ v \ d \ F) \ n \in \text{carrier } R) \wedge$
 $(\forall n. \text{deg } R \ (Vr \ K \ v) \ X \ ((Pseq1 \ R \ X \ K \ v \ d \ F) \ n) \leq an \ d) \wedge$
 $(\forall N. \exists M. \forall n \ m. M < n \wedge M < m \longrightarrow P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v \ (Vr \ K \ v) \ (an \ N))$
 $((Pseq1 \ R \ X \ K \ v \ d \ F) \ n \pm_R -_a R \ ((Pseq1 \ R \ X \ K \ v \ d \ F) \ m))))$)
apply (subgoal-tac $\forall N. \exists M. \forall n \ m. M < n \wedge M < m \longrightarrow$
 $P\text{-mod } R \ (Vr \ K \ v) \ X \ (vp \ K \ v \ (Vr \ K \ v) \ (an \ N))$
 $((Pseqh \ R \ X \ K \ v \ d \ F) \ n \pm_R -_a R \ ((Pseqh \ R \ X \ K \ v \ d \ F) \ m))))$)
apply (frule-tac $f = \lambda j. \text{snd} \ (\text{scf-d } R \ (Vr \ K \ v) \ X \ (F \ j) \ (\text{Suc } d)) \ (\text{Suc } d)$)

and $F = \text{Pseqh}_R X K v d F$ **and** $d = \text{Suc } d$ **in** $\text{monomial-}P\text{-limt}[of\ v\ R\ X]$,
 assumption+)
apply (rule allI)
apply ($\text{drule-tac } x = n$ **in** spec ,
 $\text{drule-tac } x = n$ **in** spec ,
 $\text{frule-tac } p = F\ n$ **and** $d = \text{Suc } d$ **in** $\text{PolynRg.scf-d-pol}[of\ R\ Vr\ K\ v\ X]$, assumption+ , (erule conjE) $+$)
apply ($\text{rule-tac } c = \text{scf-d } R\ (Vr\ K\ v)\ X\ (F\ n)\ (\text{Suc } d)$ **and** $j = \text{Suc } d$ **in**
 $\text{PolynRg.pol-coeff-mem}[of\ R\ Vr\ K\ v\ X]$, assumption+)
apply simp
apply (rule allI)
apply ($\text{thin-tac } \forall N. \exists M. \forall n\ m. M < n \wedge M < m \longrightarrow$
 $P\text{-mod } R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))$
 $((\text{Pseqh}_R X K v d F)\ n \pm_R -_a R ((\text{Pseqh}_R X K v d F)\ m)))$)
apply ($\text{simp only: } P\text{seqh-def hdeg-p-def, assumption, erule bexE}$)

apply ($\text{thin-tac } \forall N. \exists M. \forall n\ m. M < n \wedge M < m \longrightarrow$
 $P\text{-mod } R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))\ (F\ n \pm_R -_a R (F\ m)),$
 $\text{thin-tac } \forall N. \exists M. \forall n\ m. M < n \wedge M < m \longrightarrow$
 $P\text{-mod } R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))$
 $((\text{Pseqh}_R X K v d F)\ n \pm_R -_a R ((\text{Pseqh}_R X K v d F)\ m)))$)

apply ($\text{subgoal-tac } P\text{limit}_R X K v F\ (p \pm_R b \cdot_r R (X^{\sim R} (\text{Suc } d)))$,
 $\text{subgoal-tac } p \pm_R b \cdot_r R (X^{\sim R} (\text{Suc } d)) \in \text{carrier } R, \text{blast}$)

apply ($\text{frule } \text{PolynRg.X-mem-R}[of\ R\ Vr\ K\ v\ X]$,
 $\text{frule-tac } n = \text{Suc } d$ **in** $\text{Ring.npClose}[of\ R\ X]$, assumption+ ,
 $\text{frule } \text{Ring.ring-is-ag}[of\ R]$,
 $\text{rule aGroup.ag-pOp-closed}[of\ R]$, assumption+ ,
 $\text{rule } \text{Ring.ring-tOp-closed}$, assumption+)
apply ($\text{frule } \text{PolynRg.subring}[of\ R\ Vr\ K\ v\ X]$,
 $\text{simp add: } \text{Ring.mem-subring-mem-ring}$, assumption)

apply ($\text{simp del: npow-suc add: pol-limit-def}$,
 rule allI ,
 $\text{subgoal-tac } \forall n. F\ n = (\text{Pseql}_R X K v d F)\ n \pm_R ((\text{Pseqh}_R X K v d F)\ n)$,
 $\text{simp del: npow-suc}$,
 $\text{subgoal-tac } \forall m. (\text{Pseql}_R X K v d F)\ m \pm_R ((\text{Pseqh}_R X K v d F)\ m) \pm_R$
 $-_a R (p \pm_R (b \cdot_r R (X^{\sim R} (\text{Suc } d)))) = (\text{Pseql}_R X K v d F)\ m \pm_R -_a R p$
 $\pm_R$
 $((\text{Pseqh}_R X K v d F)\ m \pm_R -_a R (b \cdot_r R (X^{\sim R} (\text{Suc } d))))$,
 $\text{simp del: npow-suc}$)

apply ($\text{thin-tac } \forall m. (\text{Pseql}_R X K v d F)\ m \pm_R (\text{Pseqh}_R X K v d F)\ m \pm_R$
 $-_a R (p \pm_R b \cdot_r R X^{\sim R} (\text{Suc } d)) = (\text{Pseql}_R X K v d F)\ m \pm_R -_a R p \pm_R$
 $((\text{Pseqh}_R X K v d F)\ m \pm_R -_a R b \cdot_r R X^{\sim R} (\text{Suc } d)))$)
apply (erule conjE) $+$
apply ($\text{rotate-tac } -3$,

$drule-tac\ x = N$ **in** $spec$, $erule\ exE$)
apply ($rotate-tac\ 1$,
 $drule-tac\ x = N$ **in** $spec$, $erule\ exE$)
apply ($rename-tac\ d\ F\ p\ b\ N\ M1\ M2$)
apply ($subgoal-tac\ \forall m. (\max\ M1\ M2) < m \longrightarrow$
 $P-mod\ R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))$
 $((PseqI\ R\ X\ K\ v\ d\ F)\ m\ \pm_R\ -_aR\ p\ \pm_R$
 $((Pseqh\ R\ X\ K\ v\ d\ F)\ m\ \pm_R\ -_aR\ (b\ \cdot_rR\ (X^{\sim R}\ (Suc\ d))))$), $blast$)
apply ($rule\ allI$, $rule\ impI$)
apply ($rotate-tac\ -2$,
 $drule-tac\ x = m$ **in** $spec$, $simp\ del:npow-suc$)
apply ($erule\ conjE$)
apply ($rotate-tac\ -5$,
 $drule-tac\ x = m$ **in** $spec$, $simp\ del:npow-suc$)
apply ($frule-tac\ n = an\ N$ **in** $vp-apow-ideal[of\ v]$, $simp\ del:npow-suc$)
apply ($frule\ Ring.ring-is-ag[of\ R]$)
apply ($rule-tac\ I = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$ **and**
 $p = (PseqI\ R\ X\ K\ v\ d\ F)\ m\ \pm_R\ -_aR\ p$ **and**
 $q = (Pseqh\ R\ X\ K\ v\ d\ F)\ m\ \pm_R\ -_aR\ (b\ \cdot_rR\ (X^{\sim R}\ (Suc\ d)))$) **in**
 $PolynRg.P-mod-add[of\ R\ Vr\ K\ v\ X]$, $assumption+$)
apply ($rule\ aGroup.ag-pOp-closed$, $assumption+$)
apply ($simp\ add:PseqI-mem$, $rule\ aGroup.ag-mOp-closed$, $assumption+$)

apply ($rule\ aGroup.ag-pOp-closed[of\ R]$, $assumption$)
apply ($simp\ add:Pseqh-mem$, $rule\ aGroup.ag-mOp-closed$, $assumption$)
apply ($rule\ Ring.ring-tOp-closed$, $assumption$)
apply ($frule\ PolynRg.subring[of\ R\ Vr\ K\ v\ X]$)
apply ($simp\ add:Ring.mem-subring-mem-ring$)
apply ($frule\ PolynRg.X-mem-R[of\ R\ Vr\ K\ v\ X]$)
apply ($rule\ Ring.npClose$, $assumption+$)

apply ($rule\ allI$)
apply ($thin-tac\ \forall n. (PseqI\ R\ X\ K\ v\ d\ F)\ n\ \pm_R\ ((Pseqh\ R\ X\ K\ v\ d\ F)\ n) \in carrier\ R$)
apply ($erule\ conjE$)
apply ($thin-tac\ \forall n. deg\ R\ (Vr\ K\ v)\ X$
 $((PseqI\ R\ X\ K\ v\ d\ F)\ n\ \pm_R\ (Pseqh\ R\ X\ K\ v\ d\ F)\ n) \leq an\ (Suc\ d))$)
apply ($thin-tac\ \forall N. \exists M. \forall m > M. P-mod\ R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))$
 $((PseqI\ R\ X\ K\ v\ d\ F)\ m\ \pm_R\ -_aR\ p)$,
 $thin-tac\ \forall N. \exists M. \forall m > M. P-mod\ R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))$
 $((Pseqh\ R\ X\ K\ v\ d\ F)\ m\ \pm_R\ -_aR\ b\ \cdot_rR\ X^{\sim R}\ (Suc\ d))$),
 $thin-tac\ \forall n. F\ n = (PseqI\ R\ X\ K\ v\ d\ F)\ n\ \pm_R\ ((Pseqh\ R\ X\ K\ v\ d\ F)\ n)$)
apply ($drule-tac\ x = m$ **in** $spec$,
 $drule-tac\ x = m$ **in** $spec$)
apply ($subgoal-tac\ b\ \cdot_rR\ X^{\sim R}\ (Suc\ d) \in carrier\ R$)
apply ($frule\ Ring.ring-is-ag[of\ R]$)
apply ($frule-tac\ x = p$ **in** $aGroup.ag-mOp-closed[of\ R]$, $assumption+$)


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apply (subst aGroup.ag-pOp-assoc[of R], assumption+)
apply (rule aGroup.ag-mOp-closed, assumption+)
apply (rule aGroup.ag-pOp-closed, assumption+)
apply (frule-tac x1 =  $-_{aR} p$  and y1 = (Pseqh R X K v d F) m and z1 =
 $-_{aR} (b \cdot_{rR} (X \sim_R (Suc\ d)))$ ) in aGroup.ag-pOp-assoc[THEN sym, of R],
assumption+)
apply (rule aGroup.ag-mOp-closed, assumption+, simp del:npow-suc)
apply (subst aGroup.ag-pOp-assoc[THEN sym], assumption+)
apply (rule aGroup.ag-mOp-closed, assumption+,
rule aGroup.ag-pOp-closed, assumption+)
apply (subst aGroup.ag-add4-rel[of R], assumption+)
apply (rule aGroup.ag-mOp-closed, assumption+)
apply (subst aGroup.ag-p-inv[THEN sym, of R], assumption+, simp del:npow-suc)

apply (rule Ring.ring-tOp-closed, assumption+,
frule PolynRg.subring[of R Vr K v X],
simp add:Ring.mem-subring-mem-ring,
rule Ring.npClose, assumption+, simp add:PolynRg.X-mem-R)
apply (rule allI)
apply (rule-tac F = F and n = n and d = d in Pseq-decompos[of v R X],
assumption+, simp, simp)
apply (rule allI)
apply (rotate-tac -3, drule-tac x = N in spec)
apply (erule exE)
apply (subgoal-tac  $\forall n\ m. M < n \wedge M < m \longrightarrow$ 
 $P\text{-mod } R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))$ 
 $((Pseqh\ R\ X\ K\ v\ d\ F)\ n \pm_R -_{aR} ((Pseqh\ R\ X\ K\ v\ d\ F)\ m)))$ 
apply blast
apply ((rule allI)+, rule impI)
apply (rotate-tac -2,
drule-tac x = n in spec,
rotate-tac -1,
drule-tac x = m in spec,
simp)

apply (simp only: Pseqh-def)
apply (subst v-hdeg-p-mOp[of v R X], assumption+)
apply simp+

apply (frule Ring.ring-is-ag[of R])
apply (subst v-hdeg-p-pOp[of v R X], assumption+)
apply (simp, rule aGroup.ag-mOp-closed, assumption, simp, simp,
frule-tac p1 = F m in PolynRg.deg-minus-eq1 [THEN sym, of R Vr K v
X])
apply simp
apply (rotate-tac -1, frule sym,
thin-tac deg R (Vr K v) X (F m) = deg R (Vr K v) X ( $-_{aR} (F\ m)$ ), simp)
apply (rule PCauchy-hTr[of v R X], assumption+)
apply (rule-tac x = F n and y =  $-_{aR} (F\ m)$  in aGroup.ag-pOp-closed[of R],

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$assumption+$,
 $simp$, rule $aGroup.ag-mOp-closed$, $assumption+$, $simp$)
apply ($frule-tac\ p = F\ m$ **in** $PolynRg.deg-minus-eq1$ [of $R\ Vr\ K\ v\ X$],
 $simp$,
 $rule-tac\ p = F\ n$ **and** $q = -_aR\ (F\ m)$ **and** $n = Suc\ d$ **in**
 $PolynRg.polyn-deg-add4$ [of $R\ Vr\ K\ v\ X$], $assumption+$,
 $simp$, rule $aGroup.ag-mOp-closed$, $assumption$, $simp+$)

apply ($rule\ conjI$, $rule\ allI$, $rule\ PseqI-mem$, $assumption+$, $simp$)
apply ($rule\ allI$, $simp$)

apply ($thin-tac\ \forall F. (\forall n. F\ n \in carrier\ R) \wedge$
 $(\forall n. deg\ R\ (Vr\ K\ v)\ X\ (F\ n) \leq an\ d) \wedge$
 $(\forall N. \exists M. \forall n\ m. M < n \wedge M < m \longrightarrow$
 $P-mod\ R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))\ (F\ n \pm_R -_aR\ (F\ m))) \longrightarrow$
 $(\exists p \in carrier\ R. Plimit\ R\ X\ K\ v\ F\ p))$)

apply ($rule\ conjI$, $rule\ allI$)
apply ($subst\ PseqI-def$)
apply ($rule-tac\ p = F\ n$ **and** $d = d$ **in** $PolynRg.deg-ldeg-p$ [of $R\ Vr\ K\ v\ X$],
 $assumption+$)
apply $simp+$

apply ($simp\ only: PseqI-def$)
apply ($rule\ allI$)
apply ($rotate-tac\ -1$, $drule-tac\ x = N$ **in** $spec$)
apply ($erule\ exE$)
apply ($subgoal-tac\ \forall n\ m. M < n \wedge M < m \longrightarrow$
 $P-mod\ R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))$
 $(ldeg-p\ R\ (Vr\ K\ v)\ X\ d\ (F\ n) \pm_R -_aR\ (ldeg-p\ R\ (Vr\ K\ v)\ X\ d\ (F\ m)))$)
apply $blast$
apply ($(rule\ allI)+$, $rule\ impI$, $erule\ conjE$)
apply ($rotate-tac\ -3$, $drule-tac\ x = n$ **in** $spec$,
 $rotate-tac\ -1$, $drule-tac\ x = m$ **in** $spec$, $simp$)

apply ($subst\ v-ldeg-p-mOp$ [of $v\ R\ X$], $assumption+$, $simp+$)

apply ($frule\ Ring.ring-is-ag$ [of R])
apply ($subst\ v-ldeg-p-pOp$ [of $v\ R\ X$], $assumption+$, $simp$,
 $rule\ aGroup.ag-mOp-closed$, $assumption$, $simp$, $simp$,
 $frule-tac\ p = F\ m$ **in** $PolynRg.deg-minus-eq1$ [of $R\ Vr\ K\ v\ X$], $simp$)
apply $simp$
apply ($rule\ PCauchy-lTr$ [of $v\ R\ X$], $assumption+$)
apply ($rule-tac\ x = F\ n$ **and** $y = -_aR\ (F\ m)$ **in** $aGroup.ag-pOp-closed$ [of R],
 $assumption+$,
 $simp$, rule $aGroup.ag-mOp-closed$, $assumption+$, $simp$)

apply ($frule-tac\ p = F\ m$ **in** $PolynRg.deg-minus-eq1$ [of $R\ Vr\ K\ v\ X$], $simp$,
 $rule-tac\ p = F\ n$ **and** $q = -_aR\ (F\ m)$ **and** $n = Suc\ d$ **in**

$PolynRg.polyn-deg-add4[of\ R\ Vr\ K\ v\ X],\ assumption+,$
 $simp,\ rule\ aGroup.ag-mOp-closed,\ assumption,\ simp+)$
done

lemma (in Corps) $PCauchy-Plimit: [valuation\ K\ v; Complete_v\ K;$
 $PolynRg\ R\ (Vr\ K\ v)\ X; PCauchy_R\ X\ K\ v\ F] \implies$
 $\exists p \in carrier\ R. Plimit_R\ X\ K\ v\ F\ p$
by (metis $P-limitTr\ pol-Cauchy-seq-def$)

lemma (in Corps) $P-limit-mult: [valuation\ K\ v; PolynRg\ R\ (Vr\ K\ v)\ X;$
 $\forall n. F\ n \in carrier\ R; \forall n. G\ n \in carrier\ R; p1 \in carrier\ R; p2 \in carrier\ R;$
 $Plimit_R\ X\ K\ v\ F\ p1; Plimit_R\ X\ K\ v\ G\ p2] \implies$
 $Plimit_R\ X\ K\ v\ (\lambda n. (F\ n) \cdot_{rR} (G\ n))\ (p1 \cdot_{rR} p2)$
apply (frule $Vr-ring[of\ v]$,
frule $PolynRg.is-Ring$,
frule $Ring.ring-is-ag[of\ R]$)
apply (simp add: $pol-limit-def$)
apply (rule conjI)
apply (rule allI)
apply (drule-tac $x = n$ in spec,
drule-tac $x = n$ in spec)

apply (simp add: $Ring.ring-tOp-closed[of\ R]$)

apply (rule allI)
apply (rotate-tac 6,
drule-tac $x = N$ in spec,
drule-tac $x = N$ in spec)
apply (erule exE, erule exE, rename-tac $N\ M1\ M2$)
apply (subgoal-tac $\forall m. (max\ M1\ M2) < m \longrightarrow$
 $P-mod\ R\ (Vr\ K\ v)\ X\ (vp\ K\ v\ (Vr\ K\ v)\ (an\ N))$
 $((F\ m) \cdot_{rR} (G\ m) \pm_R -_{aR} (p1 \cdot_{rR} p2)))$)
apply blast

apply (rule allI, rule impI, simp, erule conjE)
apply (rotate-tac -4,
drule-tac $x = m$ in spec,
drule-tac $x = m$ in spec, simp)
apply (subgoal-tac $(F\ m) \cdot_{rR} (G\ m) \pm_R -_{aR} p1 \cdot_{rR} p2 =$
 $((F\ m) \pm_R -_{aR} p1) \cdot_{rR} (G\ m) \pm_R p1 \cdot_{rR} ((G\ m) \pm_R -_{aR} p2),\ simp)$

apply (frule-tac $n = an\ N$ in $vp-apow-ideal[of\ v]$)
apply simp
apply (rule-tac $I = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$ and
 $p = ((F\ m) \pm_R -_{aR} p1) \cdot_{rR} (G\ m)$ and $q = p1 \cdot_{rR} ((G\ m) \pm_R -_{aR} p2)$
in $PolynRg.P-mod-add[of\ R\ Vr\ K\ v\ X],$
assumption+)

apply (rule $Ring.ring-tOp-closed[of\ R],\ assumption+$)
apply (rule $aGroup.ag-pOp-closed,\ assumption$) **apply** simp

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apply (rule aGroup.ag-mOp-closed, assumption+) apply simp
apply (rule Ring.ring-tOp-closed[of R], assumption+)
apply (rule aGroup.ag-pOp-closed, assumption) apply simp
apply (rule aGroup.ag-mOp-closed, assumption+)
apply (frule Ring.whole-ideal[of Vr K v])
apply (frule-tac  $I = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$  and  $J = carrier\ (Vr\ K\ v)$  and
 $p = F\ m\ \pm_R\ -_aR\ p1$  and  $q = G\ m$  in PolynRg.P-mod-mult1[of R Vr K v X],
assumption+,
rule aGroup.ag-pOp-closed, assumption+, simp, rule aGroup.ag-mOp-closed,
assumption+) apply simp apply assumption
apply (rotate-tac 8,
drule-tac  $x = m$  in spec)
apply (case-tac  $G\ m = 0_R$ , simp add:P-mod-def)
apply (frule-tac  $p = G\ m$  in PolynRg.s-cf-expr[of R Vr K v X], assumption+,
(erule conjE)+)
thm PolynRg.P-mod-mod
apply (frule-tac  $I1 = carrier\ (Vr\ K\ v)$  and  $p1 = G\ m$  and
 $c1 = s-cf\ R\ (Vr\ K\ v)\ X\ (G\ m)$  in PolynRg.P-mod-mod[THEN sym,
of R Vr K v X], assumption+)
apply (simp,
thin-tac P-mod R (Vr K v) X (carrier (Vr K v)) (G m) =
( $\forall j \leq fst\ (s-cf\ R\ (Vr\ K\ v)\ X\ (G\ m))$ ).
snd (s-cf R (Vr K v) X (G m))  $j \in carrier\ (Vr\ K\ v)$ ))
apply (rule allI, rule impI)
apply (simp add:PolynRg.pol-coeff-mem)
apply (simp add:Ring.idealprod-whole-r[of Vr K v])

apply (cut-tac  $I = carrier\ (Vr\ K\ v)$  and  $J = vp\ K\ v\ (Vr\ K\ v)\ (an\ N)$  and
 $p = p1$  and  $q = G\ m\ \pm_R\ -_aR\ p2$  in PolynRg.P-mod-mult1[of R Vr K v X],
assumption+)
apply (simp only: Ring.whole-ideal, assumption+)
apply (rule aGroup.ag-pOp-closed, assumption+, simp, rule aGroup.ag-mOp-closed,
assumption+)
apply (frule PolynRg.s-cf-expr0[of R Vr K v X p1], assumption+)
thm PolynRg.P-mod-mod
apply (cut-tac  $I1 = carrier\ (Vr\ K\ v)$  and  $p1 = p1$  and
 $c1 = s-cf\ R\ (Vr\ K\ v)\ X\ p1$  in PolynRg.P-mod-mod[THEN sym,
of R Vr K v X], assumption+)
apply (simp add:Ring.whole-ideal, assumption+)
apply (simp, simp, simp, (erule conjE)+,
thin-tac P-mod R (Vr K v) X (carrier (Vr K v))  $p1 =$ 
( $\forall j \leq fst\ (s-cf\ R\ (Vr\ K\ v)\ X\ p1)$ ).
snd (s-cf R (Vr K v) X  $p1$ )  $j \in carrier\ (Vr\ K\ v)$ ))
apply (rule allI, rule impI)
apply (simp add:PolynRg.pol-coeff-mem, assumption)
apply (simp add:Ring.idealprod-whole-l[of Vr K v])
apply (drule-tac  $x = m$  in spec,
drule-tac  $x = m$  in spec)
apply (frule aGroup.ag-mOp-closed[of R p1], assumption,

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$(pj (Vr K v) ((Vr K v) \diamond_p t)) h) \wedge$
 $rel\text{-}prime\text{-}pols S ((Vr K v) /_r ((Vr K v) \diamond_p t)) Y$
 $(erH R (Vr K v) X S ((Vr K v) /_r ((Vr K v) \diamond_p t)) Y$
 $(pj (Vr K v) ((Vr K v) \diamond_p t)) g)$
 $(erH R (Vr K v) X S ((Vr K v) /_r ((Vr K v) \diamond_p t)) Y$
 $(pj (Vr K v) ((Vr K v) \diamond_p t)) h) \wedge$
 $P\text{-}mod R (Vr K v) X ((Vr K v) \diamond_p t) (f \pm_R -_a R (g \cdot_r R h))$
apply (frule Vr-ring[of v],
frule PolynRg.subring[of R Vr K v X],
frule vp-maximal [of v], frule PolynRg.is-Ring,
frule Ring.subring-Ring[of R Vr K v], assumption+,
frule Ring.residue-field-cd[of Vr K v vp K v], assumption+,
frule Corps.field-is-ring[of Vr K v /_r vp K v],
frule pj-Hom[of Vr K v vp K v], frule vp-ideal[of v],
simp add:Ring.maximal-ideal-ideal)
apply (frule Corps.field-is-idom[of (Vr K v) /_r (vp K v)],
frule Vr-integral[of v], simp,
frule Vr-mem-f-mem[of v t], assumption+)
apply (frule PolynRg.erH-inv[of R Vr K v X Vr K v \diamond_p t S Y g'],
assumption+, simp add:Ring.maximal-ideal-ideal,
simp add:PolynRg.is-Ring, assumption+, erule bexE, erule conjE)
apply (frule PolynRg.erH-inv[of R Vr K v X Vr K v \diamond_p t S Y h'],
assumption+, simp add:Ring.maximal-ideal-ideal,
simp add:PolynRg.is-Ring, assumption+, erule bexE, erule conjE)
apply (rename-tac g0 h0)
apply (subgoal-tac g0 $\neq \mathbf{0}_R \wedge h0 \neq \mathbf{0}_R \wedge$
 $deg R (Vr K v) X g0 \leq$
 $deg S (Vr K v /_r (Vr K v \diamond_p t)) Y (erH R (Vr K v) X S$
 $(Vr K v /_r (Vr K v \diamond_p t)) Y (pj (Vr K v) (Vr K v \diamond_p t)) g0) \wedge$
 $deg R (Vr K v) X h0 + deg S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \diamond_p t)) g0) \leq deg R (Vr K v) X f \wedge$
 $0 < deg S (Vr K v /_r (Vr K v \diamond_p t)) Y (erH R (Vr K v) X S$
 $(Vr K v /_r (Vr K v \diamond_p t)) Y (pj (Vr K v) (Vr K v \diamond_p t)) g0) \wedge$
 $0 < deg S (Vr K v /_r (Vr K v \diamond_p t)) Y (erH R (Vr K v) X S$
 $(Vr K v /_r (Vr K v \diamond_p t)) Y (pj (Vr K v) (Vr K v \diamond_p t)) h0) \wedge$
 $rel\text{-}prime\text{-}pols S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \diamond_p t)) g0)$
 $(erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \diamond_p t)) h0) \wedge$
 $P\text{-}mod R (Vr K v) X (Vr K v \diamond_p t) (f \pm_R -_a R g0 \cdot_r R h0))$
apply (thin-tac g' $\in carrier S$,
thin-tac h' $\in carrier S$,
thin-tac $0 < deg S (Vr K v /_r (Vr K v \diamond_p t)) Y g'$,
thin-tac $0 < deg S (Vr K v /_r (Vr K v \diamond_p t)) Y h'$,
thin-tac $erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \diamond_p t)) f = g' \cdot_r S h'$,
thin-tac $rel\text{-}prime\text{-}pols S (Vr K v /_r (Vr K v \diamond_p t)) Y g' h'$,

$\text{thin-tac Corps } (Vr K v /_r (Vr K v \Diamond_p t)),$
 $\text{thin-tac Ring } (Vr K v /_r (Vr K v \Diamond_p t)),$
 $\text{thin-tac vp } K v = Vr K v \Diamond_p t)$

apply *blast*
apply (*rule conjI*)
apply (*thin-tac* $0 < \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y h'$,
 $\text{thin-tac erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) f = g' \cdot_r S h'$,
 $\text{thin-tac rel-prime-pols } S (Vr K v /_r (Vr K v \Diamond_p t)) Y g' h'$,
 $\text{thin-tac erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) h0 = h'$,
 $\text{thin-tac deg } R (Vr K v) X h0 \leq \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y h'$)
apply (*rule contrapos-pp, simp+*)
apply (*simp add:PolynRg.erH-rHom-0*[*of* $R Vr K v X S$
 $Vr K v /_r (Vr K v \Diamond_p t) Y pj (Vr K v) (Vr K v \Diamond_p t)$])
apply (*rotate-tac* -3 , *drule sym*, *simp add:deg-def*)
apply (*drule aless-imp-le*[*of* $0 -\infty$],
 $\text{cut-tac minf-le-any}$ [*of* 0],
 frule ale-antisym [*of* $0 -\infty$], *simp only:ant-0*[*THEN sym*], *simp*)

apply (*rule conjI*)
apply (*thin-tac* $0 < \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y g'$,
 $\text{thin-tac erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) f = g' \cdot_r S h'$,
 $\text{thin-tac rel-prime-pols } S (Vr K v /_r (Vr K v \Diamond_p t)) Y g' h'$,
 $\text{thin-tac erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) g0 = g'$,
 $\text{thin-tac deg } R (Vr K v) X h0 \leq \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y h'$)
apply (*rule contrapos-pp, simp+*,
 $\text{simp add:PolynRg.erH-rHom-0}$ [*of* $R Vr K v X S$
 $Vr K v /_r (Vr K v \Diamond_p t) Y pj (Vr K v) (Vr K v \Diamond_p t)$])
apply (*rotate-tac* -2 , *drule sym*, *simp add:deg-def*)
apply ($\text{frule aless-imp-le}$ [*of* $0 -\infty$], *thin-tac* $0 < -\infty$,
 $\text{cut-tac minf-le-any}$ [*of* 0],
 frule ale-antisym [*of* $0 -\infty$], *simp only:ant-0*[*THEN sym*], *simp*)

apply ($\text{frule-tac } x = \text{deg } R (Vr K v) X h0$ **and**
 $y = \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y h'$ **and**
 $z = \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y g'$ **in** *aadd-le-mono*)
apply (*simp add:PolynRg.deg-mult-pols1*[*THEN sym, of* S
 $Vr K v /_r (Vr K v \Diamond_p t) Y h' g'$])
apply ($\text{frule PolynRg.is-Ring}$ [*of* $S Vr K v /_r (Vr K v \Diamond_p t) Y$],
 $\text{simp add:Ring.ring-tOp-commute}$ [*of* $S h' g'$])
apply (*rotate-tac* 11 , *drule sym*)
apply *simp*
apply ($\text{frule PolynRg.erH-rHom}$ [*of* $R Vr K v X S$
 $(Vr K v) /_r (Vr K v \Diamond_p t) Y pj (Vr K v) (Vr K v \Diamond_p t)$],
assumption+)
apply ($\text{frule PolynRg.pHom-dec-deg}$ [*of* $R Vr K v X S (Vr K v) /_r (Vr K v \Diamond_p t)$])

$Y \text{ erH } R (Vr K v) X S ((Vr K v) /_r (Vr K v \Diamond_p t))$
 $Y (pj (Vr K v) (Vr K v \Diamond_p t)) f], \text{ assumption+})$
apply (*frule-tac* $i = \text{deg } R (Vr K v) X h0 +$
 $\text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y g' \text{ in } \text{ale-trans}[\text{of } -$
 $\text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(\text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) f) \text{ deg } R (Vr K v) X f],$
 $\text{assumption+})$ **apply** *simp*
apply (*thin-tac* $\text{deg } R (Vr K v) X h0 +$
 $\text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y g' \leq \text{deg } R (Vr K v) X f,$
 $\text{thin-tac } \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(\text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) f) \leq \text{deg } R (Vr K v) X f,$
 $\text{thin-tac } 0 < \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y g',$
 $\text{thin-tac } 0 < \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y h',$
 $\text{thin-tac } \text{deg } R (Vr K v) X h0 + \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y g'$
 $\leq \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(\text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) f),$
 $\text{thin-tac } \text{rel-prime-pols } S (Vr K v /_r (Vr K v \Diamond_p t)) Y g' h')$
apply (*rotate-tac* 12, *drule sym*)
apply (*drule sym*)
apply *simp*
apply (*frule-tac* $x = g0$ and $y = h0$ in *Ring.ring-tOp-closed*[*of* R],
 $\text{assumption+})$
apply (*thin-tac* $\text{deg } R (Vr K v) X h0 \leq \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(\text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) h0),$
 $\text{thin-tac } h' = \text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) h0,$
 $\text{thin-tac } g' = \text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) g0,$
 $\text{thin-tac } \text{deg } R (Vr K v) X g0 \leq \text{deg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(\text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \Diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \Diamond_p t)) g0))$
apply (*subst PolynRg.P-mod-diff*[*THEN sym, of* $R Vr K v X Vr K v \Diamond_p t$
 $S Y f], \text{ assumption+})$ **apply** (*simp add:Ring.maximal-ideal-ideal, assumption+*)
apply (*rotate-tac* 12, *drule sym*)
apply (*subst PolynRg.erH-mult*[*of* $R Vr K v X S Vr K v /_r (Vr K v \Diamond_p t)$
 $Y], \text{ assumption+})$
done

lemma *aadd-plus-le-plus*: $\llbracket a \leq (a'::\text{ant}); b \leq b' \rrbracket \implies a + b \leq a' + b'$
by (*metis aadd-commute aadd-le-mono ale-trans*)

lemma (*in Corps*) *Hfst-PCauchy*: $\llbracket \text{valuation } K v; \text{Complete}_v K;$
 $\text{PolynRg } R (Vr K v) X; \text{PolynRg } S (Vr K v /_r (Vr K v \Diamond_p t)) Y; g0 \in \text{carrier}$
 $R;$

$h0 \in \text{carrier } R; f \in \text{carrier } R; f \neq \mathbf{0}_R; g0 \neq \mathbf{0}_R; h0 \neq \mathbf{0}_R;$
 $t \in \text{carrier } (Vr K v); vp K v = Vr K v \diamond_p t;$
 $\text{deg } R (Vr K v) X g0 \leq \text{deg } S (Vr K v /_r (Vr K v \diamond_p t)) Y (\text{erH } R (Vr K v) X S$
 $(Vr K v /_r (Vr K v \diamond_p t)) Y (pj (Vr K v) (Vr K v \diamond_p t)) g0);$
 $\text{deg } R (Vr K v) X h0 + \text{deg } S (Vr K v /_r (Vr K v \diamond_p t)) Y (\text{erH } R (Vr K v) X$
 S
 $(Vr K v /_r (Vr K v \diamond_p t)) Y (pj (Vr K v) (Vr K v \diamond_p t)) g0)$
 $\leq \text{deg } R (Vr K v) X f;$
 $0 < \text{deg } S (Vr K v /_r (Vr K v \diamond_p t)) Y (\text{erH } R (Vr K v) X S$
 $(Vr K v /_r (Vr K v \diamond_p t)) Y (pj (Vr K v) (Vr K v \diamond_p t)) g0);$
 $0 < \text{deg } S (Vr K v /_r (Vr K v \diamond_p t)) Y (\text{erH } R (Vr K v) X S$
 $(Vr K v /_r (Vr K v \diamond_p t)) Y (pj (Vr K v) (Vr K v \diamond_p t)) h0);$
 $\text{rel-prime-pols } S (Vr K v /_r (Vr K v \diamond_p t)) Y (\text{erH } R (Vr K v) X S$
 $(Vr K v /_r (Vr K v \diamond_p t)) Y (pj (Vr K v) (Vr K v \diamond_p t)) g0)$
 $(\text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \diamond_p t)) h0);$
 $\text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \diamond_p t)) f =$
 $\text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \diamond_p t)) g0 \cdot_{rS}$
 $\text{erH } R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y$
 $(pj (Vr K v) (Vr K v \diamond_p t)) h0 \rVert \implies$
 $PCauchy_R X K v \text{Hfst } K v R X t S Y f g0 h0$
apply(frule Vr-integral[of v], frule vp-ideal[of v],
frule Vr-ring, frule pj-Hom[of Vr K v vp K v], assumption+,
simp add:PolynRg.erH-mult[THEN sym, of R Vr K v X S
 $Vr K v /_r (Vr K v \diamond_p t) Y pj (Vr K v) (Vr K v \diamond_p t) g0 h0],$
frule PolynRg.P-mod-diff[THEN sym, of R Vr K v X Vr K v $\diamond_p t S Y$
 $f g0 \cdot_{rR} h0],$ assumption+,
frule PolynRg.is-Ring[of R], rule Ring.ring-tOp-closed,
assumption+)
apply (simp add:pol-Cauchy-seq-def, rule conjI)
apply (rule allI)
apply (frule-tac $t = t$ and $g = g0$ and $h = h0$ and $m = n$ in
PolynRg.P-mod-diff5-2[of R Vr K v X - S Y f],
rule Vr-integral[of v], assumption+, simp add:vp-gen-nonzero,
drule sym, simp add:vp-maximal, assumption+)

apply (subst Hfst-def)
apply (rule cart-prod-fst, assumption)
apply (rule conjI)
apply (subgoal-tac $\forall n. \text{deg } R (Vr K v) X (\text{Hfst } K v R X t S Y f g0 h0 n) \leq$
 $an (\text{deg-}n R (Vr K v) X f))$
apply blast
apply (rule allI)
apply (frule Vr-integral[of v],

$\text{frule-tac } t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = n \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-4[of } R \text{ Vr } K \text{ v } X - S \text{ Y } f], \text{ assumption+},$
 $\text{simp add:vp-gen-nonzero},$
 $\text{drule sym, simp add:vp-maximal, assumption+})$
apply ($\text{subst PolynRg.deg-an[THEN sym], (erule conjE)+, assumption+}$)
apply ($\text{simp add:Hfst-def, (erule conjE)+}$)
apply ($\text{frule-tac } i = \text{deg } R \text{ (Vr } K \text{ v) } X \text{ (fst (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ n))}$
and
 $j = \text{deg } R \text{ (Vr } K \text{ v) } X \text{ g0 and } k = \text{deg } S \text{ (Vr } K \text{ v } /_r \text{ (Vr } K \text{ v } \Diamond_p \text{ t)) } Y$
 $(\text{erH } R \text{ (Vr } K \text{ v) } X \text{ S (Vr } K \text{ v } /_r \text{ (Vr } K \text{ v } \Diamond_p \text{ t)) } Y$
 $(\text{pj (Vr } K \text{ v) (Vr } K \text{ v } \Diamond_p \text{ t)) g0) in ale-trans, assumption+},$
 $\text{frule PolynRg.nonzero-deg-pos[of } R \text{ Vr } K \text{ v } X \text{ h0}], \text{ assumption+},$
 $\text{frule-tac } x = 0 \text{ and } y = \text{deg } R \text{ (Vr } K \text{ v) } X \text{ h0 and } z = \text{deg } S$
 $(\text{Vr } K \text{ v } /_r \text{ (Vr } K \text{ v } \Diamond_p \text{ t)) } Y \text{ (erH } R \text{ (Vr } K \text{ v) } X \text{ S (Vr } K \text{ v } /_r \text{ (Vr } K \text{ v } \Diamond_p \text{ t)) } Y$
 $(\text{pj (Vr } K \text{ v) (Vr } K \text{ v } \Diamond_p \text{ t)) g0) in aadd-le-mono, simp add:aadd-0-l)}$
apply (rule allI)
apply ($\text{subgoal-tac } \forall n \text{ m. } N < n \wedge N < m \longrightarrow$
 $P\text{-mod } R \text{ (Vr } K \text{ v) } X \text{ (Vr } K \text{ v } \Diamond_p \text{ t (Vr } K \text{ v) (an } N))$
 $(\text{Hfst } K \text{ v } R \text{ X t } S \text{ Y } f \text{ g0 } h0 \text{ n } \pm_R \neg_a R \text{ (Hfst } K \text{ v } R \text{ X t } S \text{ Y } f \text{ g0 } h0 \text{ m))})$
apply blast
apply ($(\text{rule allI})+, \text{rule impI, (erule conjE)+},$
 $\text{frule Vr-integral[of v], frule vp-gen-nonzero[of v t], assumption+},$
 $\text{frule vp-maximal[of v]})$
apply ($\text{frule-tac } t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = n \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-2[of } R \text{ Vr } K \text{ v } X - S \text{ Y } f], \text{ assumption+}, \text{simp},$
 $\text{assumption+},$
 $\text{frule-tac } t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = m \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-2[of } R \text{ Vr } K \text{ v } X - S \text{ Y } f], \text{ assumption+}, \text{simp},$
 $\text{assumption+},$
 $\text{frule-tac } t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = N \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-2[of } R \text{ Vr } K \text{ v } X - S \text{ Y } f], \text{ assumption+}, \text{simp},$
 $\text{assumption+},$
 $\text{frule-tac } x = \text{Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ m in cart-prod-fst[of -}$
 $\text{carrier } R \text{ carrier } R],$
 $\text{frule-tac } x = \text{Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ N in cart-prod-fst[of -}$
 $\text{carrier } R \text{ carrier } R],$
 $\text{frule-tac } x = \text{Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ n in cart-prod-fst[of -}$
 $\text{carrier } R \text{ carrier } R],$
 $\text{thin-tac Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ n } \in \text{carrier } R \times \text{carrier } R,$
 $\text{thin-tac Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ m } \in \text{carrier } R \times \text{carrier } R)$
apply ($\text{frule PolynRg.is-Ring}$)
apply ($\text{case-tac } N = 0, \text{simp add:r-apow-def}$)
apply ($\text{rule-tac } p = \text{Hfst } K \text{ v } R \text{ X t } S \text{ Y } f \text{ g0 } h0 \text{ n } \pm_R$
 $\neg_a R \text{ (Hfst } K \text{ v } R \text{ X t } S \text{ Y } f \text{ g0 } h0 \text{ m) in}$
 $\text{PolynRg.P-mod-whole[of } R \text{ Vr } K \text{ v } X], \text{ assumption+},$
 $\text{frule Ring.ring-is-ag[of R], simp add:Hfst-def},$
 $\text{rule aGroup.ag-pOp-closed, assumption+}, \text{rule aGroup.ag-mOp-closed},$
 $\text{assumption+})$

apply (*frule-tac* $t = t$ **and** $g = g0$ **and** $h = h0$ **and** $m = N$ **and** $n = n - N$ **in**
PolynRg.P-mod-diffxxx5-3[*of* R Vr K v $X - S$ Y f], *assumption+*,
simp+, (*erule conjE*)*+*,
frule-tac $t = t$ **and** $g = g0$ **and** $h = h0$ **and** $m = N$ **and** $n = m - N$ **in**
PolynRg.P-mod-diffxxx5-3[*of* R Vr K v $X - S$ Y f], *assumption+*,
simp, (*erule conjE*)*+*,
thin-tac $P\text{-mod } R (Vr K v) X (Vr K v \diamond_p (t \neg (Vr K v) N))$ (*snd*
 $(Hpr R (Vr K v) X t S Y f g0 h0 N) \pm_R -_a R (snd (Hpr R (Vr K v) X t S Y f g0 h0$
 $n)))$),
thin-tac $P\text{-mod } R (Vr K v) X (Vr K v \diamond_p (t \neg (Vr K v) N))$ (*snd*
 $(Hpr R (Vr K v) X t S Y f g0 h0 N) \pm_R -_a R (snd (Hpr R (Vr K v) X t S Y f g0 h0$
 $m))))$)
apply (*frule* $Vr\text{-ring}[of v]$,
simp only:Ring.principal-ideal-n-pow1[*THEN sym*],
drule sym, *simp*, *frule vp-ideal*[*of v*],
simp add:Ring.ring-pow-apow,
frule-tac $n = an N$ **in** $vp\text{-apow-ideal}[of v]$, *simp*,
frule Ring.ring-is-ag[*of R*],
frule-tac $x = fst (Hpr R (Vr K v) X t S Y f g0 h0 n)$ **in**
 $aGroup.ag-mOp-closed[of R]$, *assumption*)
apply (*frule-tac* $I = vp K v (Vr K v) (an N)$ **and**
 $p = (fst (Hpr R (Vr K v) X t S Y f g0 h0 N)) \pm_R$
 $-_a R (fst (Hpr R (Vr K v) X t S Y f g0 h0 n))$ **in** *PolynRg.P-mod-minus*[*of* R
 $Vr K v X]$, *assumption+*)

apply (*rule aGroup.ag-pOp-closed*, *assumption+*,
simp add:aGroup.ag-p-inv aGroup.ag-inv-inv)
apply (*frule Ring.ring-is-ag*,
frule-tac $x = -_a R fst (Hpr R (Vr K v) X t S Y f g0 h0 N)$ **and**
 $y = fst (Hpr R (Vr K v) X t S Y f g0 h0 n)$ **in** $aGroup.ag-pOp-commute$)
apply(*rule aGroup.ag-mOp-closed*, *assumption+*, *simp*,
thin-tac $-_a R fst (Hpr R (Vr K v) X t S Y f g0 h0 N) \pm_R$
 $fst (Hpr R (Vr K v) X t S Y f g0 h0 n) = fst (Hpr R (Vr K v) X t S Y f g0 h0$
 $n) \pm_R$
 $-_a R fst (Hpr R (Vr K v) X t S Y f g0 h0 N))$)
apply (*frule-tac* $I = vp K v (Vr K v) (an N)$ **and**
 $p = fst (Hpr R (Vr K v) X t S Y f g0 h0 n) \pm_R$
 $-_a R fst (Hpr R (Vr K v) X t S Y f g0 h0 N)$ **and**
 $q = fst (Hpr R (Vr K v) X t S Y f g0 h0 N) \pm_R$
 $-_a R fst (Hpr R (Vr K v) X t S Y f g0 h0 m)$ **in** *PolynRg.P-mod-add*[*of*
 $R Vr K v X]$, *assumption+*)
apply (*rule aGroup.ag-pOp-closed*, *assumption+*,
rule aGroup.ag-mOp-closed, *assumption+*)
apply (*rule aGroup.ag-pOp-closed*, *assumption+*,
rule aGroup.ag-mOp-closed, *assumption+*)

apply (*frule-tac* $x = \text{fst } (Hpr \ R \ (Vr \ K \ v) \ X \ t \ S \ Y \ f \ g0 \ h0 \ N)$ **in**
aGroup.ag-mOp-closed, *assumption+*,
frule-tac $x = \text{fst } (Hpr \ R \ (Vr \ K \ v) \ X \ t \ S \ Y \ f \ g0 \ h0 \ m)$ **in**
aGroup.ag-mOp-closed, *assumption+*,
simp add:aGroup.pOp-assocTr43[of R] *aGroup.ag-l-inv1* *aGroup.ag-r-zero*)
apply (*simp add:Hfst-def*)
done

lemma (**in** *Corps*) *Hsnd-PCauchy*: $\llbracket valuation \ K \ v; Complete_v \ K;$
PolynRg \ R \ (Vr \ K \ v) \ X; PolynRg \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y; g0 \in carrier
R;
h0 \in carrier \ R; f \in carrier \ R; f \neq \mathbf{0}_R; g0 \neq \mathbf{0}_R; h0 \neq \mathbf{0}_R;
t \in carrier \ (Vr \ K \ v); vp \ K \ v = Vr \ K \ v \ \Diamond_p \ t;
deg \ R \ (Vr \ K \ v) \ X \ g0 \leq deg \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (erH \ R \ (Vr \ K \ v) \ X \ S
(Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ g0);
deg \ R \ (Vr \ K \ v) \ X \ h0 + deg \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (erH \ R \ (Vr \ K \ v) \ X
S
(Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ g0)
 $\leq deg \ R \ (Vr \ K \ v) \ X \ f;$
 $0 < deg \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (erH \ R \ (Vr \ K \ v) \ X \ S$
 $(Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ g0);$
 $0 < deg \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (erH \ R \ (Vr \ K \ v) \ X \ S$
 $(Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ h0);$
rel-prime-pols \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (erH \ R \ (Vr \ K \ v) \ X \ S
 $(Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \ (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ g0)$
 $(erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ h0);$
 $erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ f =$
 $erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ g0 \cdot_r S$
 $erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ h0 \rrbracket \implies$
PCauchy \ R \ X \ K \ v \ Hsnd \ K \ v \ R \ X \ t \ S \ Y \ f \ g0 \ h0

apply(*frule Vr-integral[of v]*, *frule vp-ideal[of v]*,
frule Vr-ring, *frule pj-Hom[of Vr \ K \ v \ vp \ K \ v]*, *assumption+*,
simp add:PolynRg.erH-mult[THEN sym, of R \ Vr \ K \ v \ X \ S
 $Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t) \ Y \ pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t) \ g0 \ h0]$,
frule PolynRg.P-mod-diff[THEN sym, of R \ Vr \ K \ v \ X \ Vr \ K \ v \ \Diamond_p \ t \ S \ Y
 $f \ g0 \cdot_r R \ h0]$, *assumption+*,
frule PolynRg.is-Ring[of R], *rule Ring.ring-tOp-closed*,
assumption+)
apply (*simp add:pol-Cauchy-seq-def*, *rule conjI*)
apply (*rule allI*)

apply (*frule-tac* $t = t$ **and** $g = g0$ **and** $h = h0$ **and** $m = n$ **in**
PolynRg.P-mod-diffxxx5-2[of R \ Vr \ K \ v \ X - S \ Y \ f],
rule Vr-integral[of v], *assumption+*, *simp add:vp-gen-nonzero*,

$\text{drule sym, simp add:vp-maximal, assumption+}$
apply (subst Hsnd-def)
apply (rule cart-prod-snd, assumption)
apply (rule conjI)
apply (subgoal-tac $\forall n. \text{deg } R \text{ (Vr } K \text{ v) } X \text{ (Hsnd } K \text{ v } R \text{ X } t \text{ S } Y \text{ f } g0 \text{ h0 } n) \leq$
 $\text{an (deg-n } R \text{ (Vr } K \text{ v) } X \text{ f))}$
apply blast
apply (rule allI)
apply (frule Vr-integral[of v],
 $\text{frule-tac } t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = n \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-4[of } R \text{ Vr } K \text{ v } X - S \text{ Y f], assumption+,}$
 $\text{simp add:vp-gen-nonzero,}$
 $\text{drule sym, simp add:vp-maximal, assumption+}$)
apply (subst PolynRg.deg-an[THEN sym], (erule conjE)+, assumption+)

apply (simp add:Hsnd-def)
apply (rule allI)
apply (subgoal-tac $\forall n \text{ m. } N < n \wedge N < m \longrightarrow$
 $\text{P-mod } R \text{ (Vr } K \text{ v) } X \text{ (Vr } K \text{ v } \Diamond_p t \text{ (Vr } K \text{ v) (an } N))}$
 $\text{(Hsnd } K \text{ v } R \text{ X } t \text{ S } Y \text{ f } g0 \text{ h0 } n \pm_R -_a R \text{ (Hsnd } K \text{ v } R \text{ X } t \text{ S } Y \text{ f } g0 \text{ h0 } m))$)
apply blast
apply ((rule allI)+, rule impI, (erule conjE)+)
apply (frule Vr-integral[of v], frule vp-gen-nonzero[of v t], assumption+)
apply (frule vp-maximal[of v])
apply (frule-tac $t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = n \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-2[of } R \text{ Vr } K \text{ v } X - S \text{ Y f], assumption+, simp,}$
 assumption+)
apply (frule-tac $t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = m \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-2[of } R \text{ Vr } K \text{ v } X - S \text{ Y f], assumption+, simp,}$
 assumption+,
 $\text{frule-tac } t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = N \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-2[of } R \text{ Vr } K \text{ v } X - S \text{ Y f], assumption+, simp,}$
 assumption+,
 $\text{frule-tac } x = \text{Hpr } R \text{ (Vr } K \text{ v) } X \text{ t S } Y \text{ f } g0 \text{ h0 } m \text{ in cart-prod-snd[of -}$
 $\text{carrier } R \text{ carrier } R],$
 $\text{frule-tac } x = \text{Hpr } R \text{ (Vr } K \text{ v) } X \text{ t S } Y \text{ f } g0 \text{ h0 } N \text{ in cart-prod-snd[of -}$
 $\text{carrier } R \text{ carrier } R],$
 $\text{frule-tac } x = \text{Hpr } R \text{ (Vr } K \text{ v) } X \text{ t S } Y \text{ f } g0 \text{ h0 } n \text{ in cart-prod-snd[of -}$
 $\text{carrier } R \text{ carrier } R],$
 $\text{thin-tac Hpr } R \text{ (Vr } K \text{ v) } X \text{ t S } Y \text{ f } g0 \text{ h0 } n \in \text{carrier } R \times \text{carrier } R,$
 $\text{thin-tac Hpr } R \text{ (Vr } K \text{ v) } X \text{ t S } Y \text{ f } g0 \text{ h0 } m \in \text{carrier } R \times \text{carrier } R$)
apply (frule PolynRg.is-Ring)
apply (case-tac $N = 0$, simp add:r-apow-def)
apply (rule-tac $p = \text{Hsnd } K \text{ v } R \text{ X } t \text{ S } Y \text{ f } g0 \text{ h0 } n \pm_R$
 $-_a R \text{ (Hsnd } K \text{ v } R \text{ X } t \text{ S } Y \text{ f } g0 \text{ h0 } m) \text{ in}$
 $\text{PolynRg.P-mod-whole[of } R \text{ Vr } K \text{ v } X], \text{ assumption+,}$
 $\text{frule Ring.ring-is-ag[of } R], \text{ simp add:Hsnd-def,}$

$\text{rule } aGroup.ag-pOp-closed, \text{ assumption+}, \text{ rule } aGroup.ag-mOp-closed,$
 $\text{assumption+})$

apply ($\text{frule-tac } t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = N \text{ and } n = n - N \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-3[of } R \text{ Vr } K \text{ v } X - S \text{ Y } f], \text{ assumption+})$
apply simp+
apply ($\text{erule conjE})+$
apply ($\text{frule-tac } t = t \text{ and } g = g0 \text{ and } h = h0 \text{ and } m = N \text{ and } n = m - N \text{ in}$
 $\text{PolynRg.P-mod-diffxxx5-3[of } R \text{ Vr } K \text{ v } X - S \text{ Y } f], \text{ assumption+})$
apply simp **apply** ($\text{erule conjE})+$
apply ($\text{thin-tac } P\text{-mod } R \text{ (Vr } K \text{ v) } X \text{ (Vr } K \text{ v } \diamond_p (t^{\neg(Vr \text{ } K \text{ } v) \text{ } N})) \text{ (fst}$
 $(Hpr \text{ } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } N) \pm_R -_a R \text{ (fst (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0$
 $n)))$,
 $\text{thin-tac } P\text{-mod } R \text{ (Vr } K \text{ v) } X \text{ (Vr } K \text{ v } \diamond_p (t^{\neg(Vr \text{ } K \text{ } v) \text{ } N})) \text{ (fst}$
 $(Hpr \text{ } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } N) \pm_R -_a R \text{ (fst (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0$
 $m))))$)
apply ($\text{frule Vr-ring[of v]}$)
apply ($\text{simp only:Ring.principal-ideal-n-powI[THEN sym]}$)
apply ($\text{drule sym, simp, frule vp-ideal[of v]}$)
apply ($\text{simp add:Ring.ring-pow-apow,}$
 $\text{frule-tac } n = an \text{ } N \text{ in vp-apow-ideal[of v], simp,}$
 $\text{frule Ring.ring-is-ag[of R],}$
 $\text{frule-tac } x = \text{snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } n) \text{ in}$
 $\text{aGroup.ag-mOp-closed[of R], assumption})$
apply ($\text{frule-tac } I = vp \text{ } K \text{ v (Vr } K \text{ v) (an } N) \text{ and}$
 $p = (\text{snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } N)) \pm_R$
 $-_a R (\text{snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } n)) \text{ in PolynRg.P-mod-minus[of } R$
 $\text{Vr } K \text{ v } X], \text{ assumption+})$

apply ($\text{rule } aGroup.ag-pOp-closed, \text{ assumption+},$
 $\text{simp add:aGroup.ag-p-inv aGroup.ag-inv-inv})$
apply ($\text{frule Ring.ring-is-ag,}$
 $\text{frule-tac } x = -_a R \text{ snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } N) \text{ and}$
 $y = \text{snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } n) \text{ in aGroup.ag-pOp-commute})$
apply($\text{rule } aGroup.ag-mOp-closed, \text{ assumption+}, \text{ simp,}$
 $\text{thin-tac } -_a R \text{ snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } N) \pm_R$
 $\text{snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } n) = \text{snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0$
 $n) \pm_R$
 $-_a R \text{ snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } N))$)
apply ($\text{frule-tac } I = vp \text{ } K \text{ v (Vr } K \text{ v) (an } N) \text{ and}$
 $p = \text{snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } n) \pm_R$
 $-_a R \text{ snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } N) \text{ and}$
 $q = \text{snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } N) \pm_R$
 $-_a R \text{ snd (Hpr } R \text{ (Vr } K \text{ v) } X \text{ t } S \text{ Y } f \text{ g0 } h0 \text{ } m) \text{ in PolynRg.P-mod-add[of}$
 $\text{R Vr } K \text{ v } X], \text{ assumption+})$
apply ($\text{rule } aGroup.ag-pOp-closed, \text{ assumption+},$

$\text{rule } aGroup.ag-mOp-closed, \text{ assumption+}$
apply ($\text{rule } aGroup.ag-pOp-closed, \text{ assumption+}$,
 $\text{rule } aGroup.ag-mOp-closed, \text{ assumption+}$)
apply ($\text{frule-tac } x = \text{snd } (Hpr \ R \ (Vr \ K \ v) \ X \ t \ S \ Y \ f \ g0 \ h0 \ N)$ **in**
 $aGroup.ag-mOp-closed, \text{ assumption+}$,
 $\text{frule-tac } x = \text{snd } (Hpr \ R \ (Vr \ K \ v) \ X \ t \ S \ Y \ f \ g0 \ h0 \ m)$ **in**
 $aGroup.ag-mOp-closed, \text{ assumption+}$,
 $\text{simp add:}aGroup.pOp-assocTr43[of \ R] \ aGroup.ag-l-inv1 \ aGroup.ag-r-zero$)
apply ($\text{simp add:}Hsnd-def$)
done

lemma (**in** *Corps*) *H-Plimit-f*: $[\text{valuation } K \ v; \text{Complete}_v \ K;$
 $t \in \text{carrier } (Vr \ K \ v); \text{vp } K \ v = Vr \ K \ v \diamond_p t;$
 $\text{PolynRg } R \ (Vr \ K \ v) \ X; \text{PolynRg } S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y;$
 $f \in \text{carrier } R; f \neq \mathbf{0}_R; g0 \in \text{carrier } R; h0 \in \text{carrier } R; g0 \neq \mathbf{0}_R;$
 $h0 \neq \mathbf{0}_R;$
 $0 < \text{deg } S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(\text{erH } R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \diamond_p t)) \ g0);$
 $0 < \text{deg } S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(\text{erH } R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \diamond_p t)) \ h0);$
 $\text{deg } R \ (Vr \ K \ v) \ X \ h0 +$
 $\text{deg } S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(\text{erH } R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \diamond_p t)) \ g0) \leq \text{deg } R \ (Vr \ K \ v) \ X \ f;$

 $\text{rel-prime-pols } S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(\text{erH } R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \diamond_p t)) \ g0)$
 $(\text{erH } R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \diamond_p t)) \ h0);$

 $\text{erH } R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \diamond_p t)) \ f =$
 $\text{erH } R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \diamond_p t)) \ g0 \cdot_r S$
 $\text{erH } R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \diamond_p t)) \ h0;$

 $\text{deg } R \ (Vr \ K \ v) \ X \ g0$
 $\leq \text{deg } S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(\text{erH } R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v /_r (Vr \ K \ v \diamond_p t)) \ Y$
 $(pj \ (Vr \ K \ v) \ (Vr \ K \ v \diamond_p t)) \ g0);$

 $g \in \text{carrier } R; h \in \text{carrier } R;$
 $\text{Plimit } R \ X \ K \ v \ (Hfst \ K \ v \ R \ X \ t \ S \ Y \ f \ g0 \ h0) \ g;$
 $\text{Plimit } R \ X \ K \ v \ (Hsnd \ K \ v \ R \ X \ t \ S \ Y \ f \ g0 \ h0) \ h;$
 $\text{Plimit } R \ X \ K \ v \ (\lambda n. \ (Hfst \ K \ v \ R \ X \ t \ S \ Y \ f \ g0 \ h0 \ n) \cdot_r R$

$$\begin{aligned} & (Hsnd \ K \ v \ R \ X \ t \ S \ Y \ f \ g0 \ h0 \ n)) \ (g \cdot_r R \ h)) \\ \implies & Plimit \ R \ X \ K \ v \ (\lambda n. (Hfst \ K \ v \ R \ X \ t \ S \ Y \ f \ g0 \ h0 \ n) \cdot_r R \\ & (Hsnd \ K \ v \ R \ X \ t \ S \ Y \ f \ g0 \ h0 \ n)) \ f \\ \text{apply} & (frule \ Vr\text{-integral}[of \ v], frule \ vp\text{-ideal}[of \ v], \\ & frule \ Vr\text{-ring}, frule \ pj\text{-Hom}[of \ Vr \ K \ v \ vp \ K \ v], assumption+, \\ & simp \ add:PolynRg.erH\text{-mult}[THEN \ sym, of \ R \ Vr \ K \ v \ X \ S \\ & Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t) \ Y \ pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t) \ g0 \ h0]) \\ \text{apply} & (frule \ PolynRg.P\text{-mod}\text{-diff}[of \ R \ Vr \ K \ v \ X \ Vr \ K \ v \ \Diamond_p \ t \ S \ Y \\ & f \ g0 \cdot_r R \ h0], assumption+, \\ & frule \ PolynRg.is\text{-Ring}[of \ R], rule \ Ring.ring\text{-tOp}\text{-closed}, \\ & assumption+, simp) \\ \text{apply} & (simp \ add:PolynRg.erH\text{-mult}[of \ R \ Vr \ K \ v \ X \ S \\ & Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t) \ Y \ pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t) \ g0 \ h0]) \\ \text{apply} & (frule \ PolynRg.is\text{-Ring}[of \ R]) \\ \text{apply} & (frule \ Hfst\text{-PCauchy}[of \ v \ R \ X \ S \ t \ Y \ g0 \ h0 \ f], assumption+, \\ & frule \ Hsnd\text{-PCauchy}[of \ v \ R \ X \ S \ t \ Y \ g0 \ h0 \ f], assumption+) \\ \text{apply} & (subst \ pol\text{-limit}\text{-def}) \\ \text{apply} & (rule \ conjI) \\ \text{apply} & (rule \ allI) \\ \text{apply} & (rule \ Ring.ring\text{-tOp}\text{-closed}, assumption) \\ \text{apply} & (simp \ add:pol\text{-Cauchy}\text{-seq}\text{-def}, simp \ add:pol\text{-Cauchy}\text{-seq}\text{-def}) \\ \text{apply} & (rule \ allI) \\ \text{apply} & (subgoal\text{-tac} \ \forall \ m > N. \ P\text{-mod} \ R \ (Vr \ K \ v) \ X \ (vp \ K \ v \ (Vr \ K \ v) \ (an \ N)) \\ & ((Hfst \ K \ v \ R \ X \ t \ S \ Y \ f \ g0 \ h0 \ m) \cdot_r R \ (Hsnd \ K \ v \ R \ X \ t \ S \ Y \ f \ g0 \ h0 \ m) \\ & \pm_R \neg_a R \ f)) \\ \text{apply} & blast \\ \\ \text{apply} & (rule \ allI, rule \ impI, frule \ Vr\text{-integral}[of \ v]) \\ \text{apply} & (frule\text{-tac} \ t = t \ \text{and} \ g = g0 \ \text{and} \ h = h0 \ \text{and} \ m = m - Suc \ 0 \ \text{in} \\ & PolynRg.P\text{-mod}\text{-diff}\text{xxx5-1}[of \ R \ Vr \ K \ v \ X - S \ Y], assumption+, \\ & simp \ add:vp\text{-gen}\text{-nonzero}[of \ v], \\ & frule \ vp\text{-maximal}[of \ v], simp, assumption+) \\ \text{apply} & ((erule \ conjE)+, simp \ del:npow\text{-suc} \ Hpr\text{-Suc}) \\ \text{apply} & (frule \ Ring.ring\text{-is}\text{-ag}[of \ R]) \\ \text{apply} & (thin\text{-tac} \ 0 < deg \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \\ & (erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \\ & (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ g0), \\ & thin\text{-tac} \ 0 < deg \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \\ & (erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \\ & (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ h0), \\ & thin\text{-tac} \ rel\text{-prime}\text{-pols} \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \\ & (erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \\ & (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ g0) \\ & (erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \\ & (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ h0), \\ & thin\text{-tac} \ P\text{-mod} \ R \ (Vr \ K \ v) \ X \ (Vr \ K \ v \ \Diamond_p \ t) \ (f \pm_R \neg_a R \ g0 \cdot_r R \ h0), \\ & thin\text{-tac} \ erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y \\ & (pj \ (Vr \ K \ v) \ (Vr \ K \ v \ \Diamond_p \ t)) \ f = \\ & (erH \ R \ (Vr \ K \ v) \ X \ S \ (Vr \ K \ v \ /_r \ (Vr \ K \ v \ \Diamond_p \ t)) \ Y
\end{aligned}$$

$$\begin{aligned}
& (pj (Vr K v) (Vr K v \diamond_p t)) g0) \cdot_r S \\
& (erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y \\
& (pj (Vr K v) (Vr K v \diamond_p t)) h0), \\
thin-tac \ deg R (Vr K v) X g0 \leq \ & deg S (Vr K v /_r (Vr K v \diamond_p t)) Y \\
& (erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y \\
& (pj (Vr K v) (Vr K v \diamond_p t)) g0), \\
thin-tac \ deg R (Vr K v) X h0 + \ & deg S (Vr K v /_r (Vr K v \diamond_p t)) Y \\
& (erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y \\
& (pj (Vr K v) (Vr K v \diamond_p t)) g0) \leq \ & deg R (Vr K v) X f, \\
thin-tac \ erH R (Vr K v) X S (Vr K v /_r \ & (Vr K v \diamond_p t)) Y \\
& (pj (Vr K v) (Vr K v \diamond_p t)) (fst (Hpr R (Vr K v) X t S Y f g0 h0 m)) = \\
erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) \ & Y \\
& (pj (Vr K v) (Vr K v \diamond_p t)) g0, \\
thin-tac \ erH R (Vr K v) X S (Vr K v /_r \ & (Vr K v \diamond_p t)) Y \\
& (pj (Vr K v) (Vr K v \diamond_p t)) (snd (Hpr R (Vr K v) X t S Y f g0 h0 m)) = \\
erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) \ & Y \\
& (pj (Vr K v) (Vr K v \diamond_p t)) h0, \\
thin-tac \ deg R (Vr K v) X (fst (Hpr R (Vr K v) X t S Y f g0 h0 m)) \\
\leq \ deg S (Vr K v /_r (Vr K v \diamond_p t)) \ & Y \\
& (erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y \\
& (pj (Vr K v) (Vr K v \diamond_p t)) g0), \\
thin-tac \ P-mod R (Vr K v) X (Vr K v \diamond_p \ & (t \neg (Vr K v) m)) \\
& (fst (Hpr R (Vr K v) X t S Y f g0 h0 (m - Suc 0)) \pm_R \\
& -_a R (fst (Hpr R (Vr K v) X t S Y f g0 h0 m))), \\
thin-tac \ deg R (Vr K v) X (snd (Hpr R (Vr K v) X t S Y f g0 h0 m)) + \\
deg S (Vr K v /_r (Vr K v \diamond_p t)) \ & Y \\
& (erH R (Vr K v) X S (Vr K v /_r (Vr K v \diamond_p t)) Y \\
& (pj (Vr K v) (Vr K v \diamond_p t)) g0) \leq \ deg R (Vr K v) X f, \\
thin-tac \ P-mod R (Vr K v) X (Vr K v \diamond_p \ & (t \neg (Vr K v) m)) \\
& (snd (Hpr R (Vr K v) X t S Y f g0 h0 (m - Suc 0)) \pm_R \\
& -_a R (snd (Hpr R (Vr K v) X t S Y f g0 h0 m)))) \\
apply (case-tac N = 0, simp add:r-apow-def) \\
apply (rule-tac p = (Hfst K v R X t S Y f g0 h0 m) \cdot_r R (Hsnd K v R X t S Y f g0 h0 m) \\
m) \\
\pm_R -_a R f \textbf{ in } PolynRg.P-mod-whole[of R Vr K v X], assumption+ \\
apply (simp add:Hfst-def Hsnd-def) \\
apply (frule-tac x = Hpr R (Vr K v) X t S Y f g0 h0 m \textbf{ in } cart-prod-fst[of - carrier R carrier R]) \\
apply (frule-tac x = Hpr R (Vr K v) X t S Y f g0 h0 m \textbf{ in } cart-prod-snd[of - carrier R carrier R]) \\
apply (frule Ring.ring-is-ag[of R], \\
rule aGroup.ag-pOp-closed, assumption) \\
apply (rule Ring.ring-tOp-closed, assumption+) \\
apply (rule aGroup.ag-mOp-closed, assumption+) \\
apply (frule-tac g = f \pm_R -_a R (fst (Hpr R (Vr K v) X t S Y f g0 h0 m)) \cdot_r R \\
(snd (Hpr R (Vr K v) X t S Y f g0 h0 m)) \textbf{ and } m = N - Suc 0 \textbf{ and } n = m
\end{aligned}$$

in
PolynRg.P-mod-n-m[of R Vr K v X], *assumption+*)
apply (*frule-tac* $x = Hpr$ R (Vr K v) X t S Y f $g0$ $h0$ m **in** *cart-prod-fst*[of - *carrier* R *carrier* R],
frule-tac $x = Hpr$ R (Vr K v) X t S Y f $g0$ $h0$ m **in** *cart-prod-snd*[of - *carrier* R *carrier* R])
apply (*rule* *aGroup.ag-pOp-closed*, *assumption+*, *rule* *aGroup.ag-mOp-closed*,
assumption)
apply (*rule* *Ring.ring-tOp-closed*, *assumption+*)
apply (*subst* *Suc-le-mono*[*THEN sym*], *simp*)
apply *assumption*
apply (*simp* *del:npow-suc*)
apply (*simp* *only:Ring.principal-ideal-n-pow1*[*THEN sym*, of Vr K v])
apply (*cut-tac* $n = N$ **in** *an-neq-inf*)
apply (*subgoal-tac* *an* $N \neq 0$)
apply (*subst* *r-apow-def*, *simp*) **apply** (*simp* *add:na-an*)
apply (*frule* *Ring.principal-ideal*[of Vr K v t], *assumption*)
apply (*frule-tac* $I = Vr$ K v $\diamond_p t$ **and** $n = N$ **in** *Ring.ideal-pow-ideal*[of Vr K v],
assumption+)
apply (*frule-tac* $x = Hpr$ R (Vr K v) X t S Y f $g0$ $h0$ m **in** *cart-prod-fst*[of - *carrier* R *carrier* R],
frule-tac $x = Hpr$ R (Vr K v) X t S Y f $g0$ $h0$ m **in** *cart-prod-snd*[of - *carrier* R *carrier* R])
apply (*thin-tac* *PCauchy* R X K v *Hfst* K v R X t S Y f $g0$ $h0$,
thin-tac *PCauchy* R X K v *Hsnd* K v R X t S Y f $g0$ $h0$)
apply (*simp* *add:Hfst-def* *Hsnd-def*)
apply (*frule-tac* $x = fst$ (Hpr R (Vr K v) X t S Y f $g0$ $h0$ m) **and** $y = snd$ (Hpr R (Vr K v) X t S Y f $g0$ $h0$ m) **in** *Ring.ring-tOp-closed*[of R], *assumption+*)
apply (*frule-tac* $x = (fst$ (Hpr R (Vr K v) X t S Y f $g0$ $h0$ m)) $\cdot_r R$ (snd (Hpr R (Vr K v) X t S Y f $g0$ $h0$ m)) **in** *aGroup.ag-mOp-closed*[of R], *assumption+*)
apply (*frule-tac* $I = Vr$ K v $\diamond_p t$ $\diamond(Vr$ K $v)$ N **and**
 $p = f \pm_R -_a R$ (fst (Hpr R (Vr K v) X t S Y f $g0$ $h0$ m)) $\cdot_r R$
 $(snd$ (Hpr R (Vr K v) X t S Y f $g0$ $h0$ m)) **in**
PolynRg.P-mod-minus[of R Vr K v X], *assumption+*)
apply (*frule* *Ring.ring-is-ag*[of R])
apply (*rule* *aGroup.ag-pOp-closed*, *assumption+*)
apply (*simp* *add:aGroup.ag-p-inv*, *simp* *add:aGroup.ag-inv-inv*,
frule *aGroup.ag-mOp-closed*[of R f], *assumption+*)
apply (*simp* *add:aGroup.ag-pOp-commute*[of R $-_a R$ f])
apply (*subst* *an-0*[*THEN sym*])
apply (*subst* *aneq-natneq*[of - 0], *thin-tac* *an* $N \neq \infty$, *simp*)
done
theorem (**in** *Corps*) *Hensel*: $\llbracket valuation$ K v ; *Complete* $_v$ K ;
PolynRg R (Vr K v) X ; *PolynRg* S ((Vr K v) $/_r$ (vp K v)) Y ;

$f \in \text{carrier } R; f \neq \mathbf{0}_R; g' \in \text{carrier } S; h' \in \text{carrier } S;$
 $0 < \deg S ((Vr K v) /_r (vp K v)) Y g';$
 $0 < \deg S ((Vr K v) /_r (vp K v)) Y h';$
 $((erH R (Vr K v) X S ((Vr K v) /_r (vp K v)) Y$
 $\quad (pj (Vr K v) (vp K v))) f) = g' \cdot_r S h';$
 $\text{rel-prime-pols } S ((Vr K v) /_r (vp K v)) Y g' h' \rrbracket \implies$
 $\exists g h. g \in \text{carrier } R \wedge h \in \text{carrier } R \wedge$
 $\deg R (Vr K v) X g \leq \deg S ((Vr K v) /_r (vp K v)) Y g' \wedge$
 $f = g \cdot_r R h$
apply (frule PolynRg.is-Ring[of $R Vr K v X$],
frule PolynRg.is-Ring[of $S Vr K v /_r vp K v Y$],
frule vp-gen-t[of v], erule bexE,
frule-tac $t = t$ **in** vp-gen-nonzero[of v], assumption)
apply (frule-tac $t = t$ **in** Hensel-starter[of $v R X S Y - f g' h'$], assumption+)
apply ((erule exE)+, (erule conjE)+, rename-tac $g0 h0$)
apply (frule Vr-ring[of v], frule Vr-integral[of v])
apply (rotate-tac 22, drule sym, drule sym, simp)
apply (frule vp-maximal[of v], simp)
apply (frule-tac $mx = Vr K v \Diamond_p t$ **in** Ring.residue-field-cd[of $Vr K v$],
assumption)
apply (frule-tac $mx = Vr K v \Diamond_p t$ **in** Ring.maximal-ideal-ideal[of $Vr K v$],
assumption)
apply (frule-tac $I = Vr K v \Diamond_p t$ **in** Ring.qring-ring[of $Vr K v$],
assumption+)
apply (frule-tac $B = Vr K v /_r (Vr K v \Diamond_p t)$ **and** $h = pj (Vr K v) (Vr K v \Diamond_p$
 $t)$ **in** PolynRg.erH-rHom[of $R Vr K v X S - Y$], assumption+)
apply (simp add: pj-Hom)
apply (frule-tac $?g0.0 = g0$ **and** $?h0.0 = h0$ **in** Hfst-PCauchy[of $v R X S - Y -$
 $-$
 f], assumption+)
apply (frule-tac $?g0.0 = g0$ **and** $?h0.0 = h0$ **in** Hsnd-PCauchy[of $v R X S - Y -$
 $-$
 f], assumption+)
apply (frule-tac $F = \lambda j. Hfst_{K v R X t S Y f g0 h0 j}$ **in** PCauchy-Plimit[of $v R X$]
, assumption+)
apply (frule-tac $F = \lambda j. Hsnd_{K v R X t S Y f g0 h0 j}$ **in** PCauchy-Plimit[of $v R$
 X]
, assumption+)
apply ((erule bexE)+, rename-tac $g0 h0 g h$)

apply (frule-tac $F = \lambda j. Hfst_{K v R X t S Y f g0 h0 j}$ **and**
 $G = \lambda j. Hsnd_{K v R X t S Y f g0 h0 j}$ **and** $?p1.0 = g$ **and** $?p2.0 = h$
in P-limit-mult[of $v R X$], assumption+, rule allI)

apply (simp add: pol-Cauchy-seq-def)
apply (simp add: pol-Cauchy-seq-def, assumption+)
apply (frule-tac $t = t$ **and** $?g0.0 = g0$ **and** $?h0.0 = h0$ **and** $g = g$ **and** $h = h$
in
H-Plimit-f[of $v - R X S Y f$], assumption+)

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apply (frule-tac  $F = \lambda n. (Hfst\ K\ v\ R\ X\ t\ S\ Y\ f\ g0\ h0\ n) \cdot_r R\ (Hsnd\ K\ v\ R\ X\ t\ S\ Y\ f\ g0\ h0\ n)$  and  $?p1.0 = g \cdot_r R\ h$  and  $d = na\ (deg\ R\ (Vr\ K\ v)\ X\ g0 + deg\ R\ (Vr\ K\ v)\ X\ f)$  in  $P\text{-limit-unique}[of\ v\ R\ X\ -\ -\ f],\ assumption+$ )
apply (rule allI)
apply (frule PolynRg.is-Ring[of  $R$ ],
      rule Ring.ring-tOp-closed, assumption)
apply (simp add:pol-Cauchy-seq-def)
apply (simp add:pol-Cauchy-seq-def)
apply (rule allI)
apply (thin-tac Plimit  $R\ X\ K\ v\ (Hfst\ K\ v\ R\ X\ t\ S\ Y\ f\ g0\ h0)\ g,$ 
      thin-tac Plimit  $R\ X\ K\ v\ (Hsnd\ K\ v\ R\ X\ t\ S\ Y\ f\ g0\ h0)\ h,$ 
      thin-tac Plimit  $R\ X\ K\ v\ (\lambda n. (Hfst\ K\ v\ R\ X\ t\ S\ Y\ f\ g0\ h0\ n) \cdot_r R\ (Hsnd\ K\ v\ R\ X\ t\ S\ Y\ f\ g0\ h0\ n))\ (g \cdot_r R\ h),$ 
      thin-tac Plimit  $R\ X\ K\ v\ (\lambda n. (Hfst\ K\ v\ R\ X\ t\ S\ Y\ f\ g0\ h0\ n) \cdot_r R\ (Hsnd\ K\ v\ R\ X\ t\ S\ Y\ f\ g0\ h0\ n))\ f)$ 
apply (subst PolynRg.deg-mult-pols1[of  $R\ Vr\ K\ v\ X$ ], assumption+)

apply (simp add:pol-Cauchy-seq-def, simp add:pol-Cauchy-seq-def,
      thin-tac  $g' = erH\ R\ (Vr\ K\ v)\ X\ S\ (Vr\ K\ v\ /_r\ (Vr\ K\ v\ \diamond_p\ t))\ Y\ (pj\ (Vr\ K\ v)\ (Vr\ K\ v\ \diamond_p\ t))\ g0,$ 
      thin-tac  $h' = erH\ R\ (Vr\ K\ v)\ X\ S\ (Vr\ K\ v\ /_r\ (Vr\ K\ v\ \diamond_p\ t))\ Y\ (pj\ (Vr\ K\ v)\ (Vr\ K\ v\ \diamond_p\ t))\ h0)$ 
apply (subst PolynRg.deg-mult-pols1[THEN sym, of  $R\ Vr\ K\ v\ X$ ], assumption+)
apply (simp add:pol-Cauchy-seq-def, simp add:pol-Cauchy-seq-def)

apply (frule-tac  $p1 = g0$  in  $PolynRg.deg-mult-pols1[THEN\ sym,\ of\ R\ Vr\ K\ v\ X - f],\ assumption+, simp)$ 
apply (frule-tac  $x = g0$  in  $Ring.ring-tOp-closed[of\ R - f],\ assumption+$ )
apply (frule-tac  $p = g0 \cdot_r R\ f$  in  $PolynRg.nonzero-deg-pos[of\ R\ Vr\ K\ v\ X],\ assumption+$ )
apply (frule-tac  $p = g0 \cdot_r R\ f$  in  $PolynRg.deg-in-aug-minf[of\ R\ Vr\ K\ v\ X],\ assumption+, simp\ add:aug-minf-def,$ 
      simp add:PolynRg.polyn-ring-integral[of  $R\ Vr\ K\ v\ X$ ],
      simp add:Idomain.idom-tOp-nonzeros[of  $R - f$ ],
      frule-tac  $p = g0 \cdot_r R\ f$  in  $PolynRg.deg-noninf[of\ R\ Vr\ K\ v\ X],\ assumption+$ )
apply (simp add:an-na)
apply (subst PolynRg.deg-mult-pols1[of  $R\ Vr\ K\ v\ X$ ], assumption+,
      simp add:pol-Cauchy-seq-def, simp add:pol-Cauchy-seq-def)
apply (frule-tac  $t = t$  and  $g = g0$  and  $h = h0$  and  $m = n$  in
       $PolynRg.P\text{-mod-diffxxx5-4}[of\ R\ Vr\ K\ v\ X - S\ Y\ f],\ assumption+$ )
apply (erule conjE)

apply (frule-tac  $a = deg\ R\ (Vr\ K\ v)\ X\ (fst\ (Hpr\ R\ (Vr\ K\ v)\ X\ t\ S\ Y\ f\ g0\ h0\ n))$ 
and  $a' = deg\ R\ (Vr\ K\ v)\ X\ g0$  and  $b = deg\ R\ (Vr\ K\ v)\ X\ (snd\ (Hpr\ R\ (Vr\ K\ v)\ X\ t\ S\ Y\ f\ g0\ h0\ n))$ 
and  $b' = deg\ R\ (Vr\ K\ v)\ X\ f$  in  $aadd-plus-le-plus,\ assumption+$ )
apply simp
apply (simp add:Hfst-def Hsnd-def)

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apply (rule Ring.ring-tOp-closed, assumption+)
apply (rule rotate-tac -1, drule sym)
apply (frule-tac  $F = \lambda j. Hfst_{K\ v\ R\ X\ t\ S\ Y\ f\ g0\ h0\ j}$  and  $p = g$  and  $ad = deg\ S$ 
  ( $Vr\ K\ v\ /\_r\ (Vr\ K\ v\ \Diamond_p\ t)$ )  $Y$  (erH  $R\ (Vr\ K\ v)\ X\ S\ (Vr\ K\ v\ /\_r\ (Vr\ K\ v\ \Diamond_p\ t))\ Y$ 
  ( $pj\ (Vr\ K\ v)\ (Vr\ K\ v\ \Diamond_p\ t)\ g0$ ) in Plimit-deg1[of v R X], assumption+,
  simp add:pol-Cauchy-seq-def)
apply (rule allI)
apply (frule-tac  $t = t$  and  $g = g0$  and  $h = h0$  and  $m = n$  in
  PolynRg.P-mod-diffxxx5-4[of R Vr K v X - S Y f], assumption+,
  erule conjE)
apply (frule-tac  $i = deg\ R\ (Vr\ K\ v)\ X\ (fst\ (Hpr\ R\ (Vr\ K\ v)\ X\ t\ S\ Y\ f\ g0\ h0\ n))$ 
and
 $j = deg\ R\ (Vr\ K\ v)\ X\ g0$  and  $k = deg\ S\ (Vr\ K\ v\ /\_r\ (Vr\ K\ v\ \Diamond_p\ t))\ Y$ 
  (erH  $R\ (Vr\ K\ v)\ X\ S\ (Vr\ K\ v\ /\_r\ (Vr\ K\ v\ \Diamond_p\ t))\ Y$ 
  ( $pj\ (Vr\ K\ v)\ (Vr\ K\ v\ \Diamond_p\ t)\ g0$ ) in ale-trans, assumption+)
apply (subst Hfst-def, assumption+, blast)
done

end

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