

The Twelfold Way

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Abstract

This entry provides all cardinality theorems of the Twelfold Way. The Twelfold Way [1, 5, 6] systematically classifies twelve related combinatorial problems concerning two finite sets, which include counting permutations, combinations, multisets, set partitions and number partitions. This development builds upon the existing formal developments [2, 3, 4] with cardinality theorems for those structures. It provides twelve bijections from the various structures to different equivalence classes on finite functions, and hence, proves cardinality formulae for these equivalence classes on finite functions.

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1 Preliminaries

theory *Preliminaries*

imports

Main

HOL-Library.Multiset

HOL-Library.FuncSet

HOL-Combinatorics.Permutations

HOL-ex.Birthday-Paradox

Card-Partitions.Card-Partitions

Bell-Numbers-Spivey.Bell-Numbers

Card-Multisets.Card-Multisets

Card-Number-Partitions.Card-Number-Partitions

begin

1.1 Additions to Finite Set Theory

lemma *subset-with-given-card-exists*:

assumes $n \leq \text{card } A$

shows $\exists B \subseteq A. \text{card } B = n$

using *assms* **proof** (*induct n*)

case 0

then show *?case* **by** *auto*

next

case (*Suc n*)

from *this* **obtain** *B* **where** $B \subseteq A$ $\text{card } B = n$ **by** *auto*

from *this* $\langle B \subseteq A \rangle \langle \text{card } B = n \rangle$ **have** $\text{card } B < \text{card } A$

using *Suc.prem*s **by** *linarith*

from $\langle \text{Suc } n \leq \text{card } A \rangle$ *card.infinite* **have** *finite A* **by** *force*

from *this* $\langle B \subseteq A \rangle$ *finite-subset* **have** *finite B* **by** *blast*

from $\langle \text{card } B < \text{card } A \rangle \langle B \subseteq A \rangle$ **obtain** *a* **where** $a \in A$ $a \notin B$

by (*metis less-irrefl subsetI subset-antisym*)

have $\text{insert } a \ B \subseteq A$ $\text{card } (\text{insert } a \ B) = \text{Suc } n$

using $\langle \text{finite } B \rangle \langle a \in A \rangle \langle a \notin B \rangle \langle B \subseteq A \rangle \langle \text{card } B = n \rangle$ **by** *auto*

then show *?case* **by** *blast*

qed

1.2 Additions to Equiv Relation Theory

lemmas *univ-commute'* = *univ-commute*[*unfolded Equiv-Relations.proj-def*]

lemma *univ-predicate-impl-forall*:

```

assumes equiv A R
assumes P respects R
assumes  $X \in A // R$ 
assumes univ P X
shows  $\forall x \in X. P x$ 
proof –
  from assms(1,3) obtain  $x$  where  $x \in X$ 
  by (metis equiv-class-self quotientE)
  from  $\langle x \in X \rangle$  assms(1,3) have  $X = R `` \{x\}$ 
  by (metis Image-singleton-iff equiv-class-eq quotientE)
  from assms(1,2,4) this show ?thesis
  using equiv-class-eq-iff univ-commute' by fastforce
qed

```

lemma *univ-preserves-predicate*:

```

assumes equiv A r
assumes P respects r
shows  $\{x \in A. P x\} // r = \{X \in A // r. \text{univ } P X\}$ 
proof
  show  $\{x \in A. P x\} // r \subseteq \{X \in A // r. \text{univ } P X\}$ 
  proof
    fix  $X$ 
    assume  $X \in \{x \in A. P x\} // r$ 
    from this obtain  $x$  where  $x \in \{x \in A. P x\}$  and  $X = r `` \{x\}$ 
    using quotientE by blast
    have  $X \in A // r$ 
    using  $\langle X = r `` \{x\} \rangle \langle x \in \{x \in A. P x\} \rangle$ 
    by (auto intro: quotientI)
    moreover have univ P X
    using  $\langle X = r `` \{x\} \rangle \langle x \in \{x \in A. P x\} \rangle$  assms
    by (simp add: proj-def[symmetric] univ-commute)
    ultimately show  $X \in \{X \in A // r. \text{univ } P X\}$  by auto
  qed
next
  show  $\{X \in A // r. \text{univ } P X\} \subseteq \{x \in A. P x\} // r$ 
  proof
    fix  $X$ 
    assume  $X \in \{X \in A // r. \text{univ } P X\}$ 
    from this have  $X \in A // r$  and univ P X by auto
    from  $\langle X \in A // r \rangle$  obtain  $x$  where  $x \in A$  and  $X = r `` \{x\}$ 
    using quotientE by blast
    have  $x \in \{x \in A. P x\}$ 
    using  $\langle x \in A \rangle \langle X = r `` \{x\} \rangle \langle \text{univ } P X \rangle$  assms
    by (simp add: proj-def[symmetric] univ-commute)
    from this show  $X \in \{x \in A. P x\} // r$ 
    using  $\langle X = r `` \{x\} \rangle$  by (auto intro: quotientI)
  qed

```

qed

lemma *Union-quotient-restricted:*

assumes *equiv A r*

assumes *P respects r*

shows $\bigcup(\{x \in A. P\ x\} // r) = \{x \in A. P\ x\}$

proof

show $\bigcup(\{x \in A. P\ x\} // r) \subseteq \{x \in A. P\ x\}$

proof

fix *x*

assume $x \in \bigcup(\{x \in A. P\ x\} // r)$

from *this* obtain *X* where $x \in X$ and $X \in \{x \in A. P\ x\} // r$ by *blast*

from *this* obtain *x'* where $X = r\ \{\!\!'\{x'\}\!\!'\}$ and $x' \in \{x \in A. P\ x\}$

using *quotientE* by *blast*

from *this* $\langle x \in X \rangle$ have $x \in A$

using $\langle \text{equiv } A\ r \rangle$ by (*simp add: equiv-class-eq-iff*)

moreover from $\langle X = r\ \{\!\!'\{x'\}\!\!'\} \rangle \langle x \in X \rangle \langle x' \in \{x \in A. P\ x\} \rangle$ have *P x*

using $\langle P\ \text{respects } r \rangle$ *congruentD* by *fastforce*

ultimately show $x \in \{x \in A. P\ x\}$ by *auto*

qed

next

show $\{x \in A. P\ x\} \subseteq \bigcup(\{x \in A. P\ x\} // r)$

proof

fix *x*

assume $x \in \{x \in A. P\ x\}$

from *this* have $x \in r\ \{\!\!'\{x}\!\!'\}$

using $\langle \text{equiv } A\ r \rangle$ *equiv-class-self* by *fastforce*

from $\langle x \in \{x \in A. P\ x\} \rangle$ have $r\ \{\!\!'\{x}\!\!'\} \in \{x \in A. P\ x\} // r$

by (*auto intro: quotientI*)

from *this* $\langle x \in r\ \{\!\!'\{x}\!\!'\} \rangle$ show $x \in \bigcup(\{x \in A. P\ x\} // r)$ by *auto*

qed

qed

lemma *finite-equiv-implies-finite-carrier:*

assumes *equiv A R*

assumes *finite (A // R)*

assumes $\forall X \in A // R. \text{finite } X$

shows *finite A*

proof –

from $\langle \text{equiv } A\ R \rangle$ have $A = \bigcup(A // R)$

by (*simp add: Union-quotient*)

from *this* $\langle \text{finite } (A // R) \rangle \langle \forall X \in A // R. \text{finite } X \rangle$ show *finite A*

using *finite-Union* by *fastforce*

qed

lemma *finite-quotient-iff:*

assumes *equiv A R*

shows $\text{finite } A \longleftrightarrow (\text{finite } (A // R) \wedge (\forall X \in A // R. \text{finite } X))$

using *assms* by (*meson equiv-type finite-equiv-class finite-equiv-implies-finite-carrier*)

finite-quotient)

1.2.1 Counting Sets by Splitting into Equivalence Classes

lemma *card-equiv-class-restricted*:

assumes *finite* $\{x \in A. P\ x\}$
assumes *equiv* $A\ R$
assumes P *respects* R
shows $\text{card } \{x \in A. P\ x\} = \text{sum } \text{card } (\{x \in A. P\ x\} // R)$
proof –
have $\text{card } \{x \in A. P\ x\} = \text{card } (\bigcup (\{x \in A. P\ x\} // R))$
using $\langle \text{equiv } A\ R \rangle \langle P \text{ respects } R \rangle$ **by** (*simp add: Union-quotient-restricted*)
also have $\text{card } (\bigcup (\{x \in A. P\ x\} // R)) = (\sum C \in \{x \in A. P\ x\} // R. \text{card } C)$
proof –
from $\langle \text{finite } \{x \in A. P\ x\} \rangle$ **have** $\text{finite } (\{x \in A. P\ x\} // R)$
using $\langle \text{equiv } A\ R \rangle$ **by** (*metis finite-imageI proj-image*)
moreover from $\langle \text{finite } \{x \in A. P\ x\} \rangle$ **have** $\forall C \in \{x \in A. P\ x\} // R. \text{finite } C$
using $\langle \text{equiv } A\ R \rangle \langle P \text{ respects } R \rangle$ *Union-quotient-restricted*
Union-upper finite-subset **by** *fastforce*
moreover have $\forall C1 \in \{x \in A. P\ x\} // R. \forall C2 \in \{x \in A. P\ x\} // R. C1 \neq C2 \longrightarrow C1 \cap C2 = \{\}$
using $\langle \text{equiv } A\ R \rangle$ *quotient-disj*
by (*metis (no-types, lifting) mem-Collect-eq quotientE quotientI*)
ultimately show *?thesis*
by (*subst card-Union-disjoint*) (*auto simp: pairwise-def disjnt-def*)
qed
finally show *?thesis* .
qed

lemma *card-equiv-class-restricted-same-size*:

assumes *equiv* $A\ R$
assumes P *respects* R
assumes $\bigwedge F. F \in \{x \in A. P\ x\} // R \implies \text{card } F = k$
shows $\text{card } \{x \in A. P\ x\} = k * \text{card } (\{x \in A. P\ x\} // R)$
proof *cases*
assume *finite* $\{x \in A. P\ x\}$
have $\text{card } \{x \in A. P\ x\} = \text{sum } \text{card } (\{x \in A. P\ x\} // R)$
using $\langle \text{finite } \{x \in A. P\ x\} \rangle \langle \text{equiv } A\ R \rangle \langle P \text{ respects } R \rangle$
by (*simp add: card-equiv-class-restricted*)
also have $\text{sum } \text{card } (\{x \in A. P\ x\} // R) = k * \text{card } (\{x \in A. P\ x\} // R)$
by (*simp add: $\langle \bigwedge F. F \in \{x \in A. P\ x\} // R \implies \text{card } F = k \rangle$*)
finally show *?thesis* .
next
assume *infinite* $\{x \in A. P\ x\}$
from this have *infinite* $(\bigcup (\{a \in A. P\ a\} // R))$
using $\langle \text{equiv } A\ R \rangle \langle P \text{ respects } R \rangle$ **by** (*simp add: Union-quotient-restricted*)
from this have *infinite* $(\{x \in A. P\ x\} // R) \vee (\exists X \in \{x \in A. P\ x\} // R. \text{infinite } X)$
by *auto*

```

from this show ?thesis
proof
  assume infinite  $\{x \in A. P\ x\} // R$ 
  from this  $\langle \text{infinite } \{x \in A. P\ x\} \rangle$  show ?thesis by simp
next
  assume  $\exists X \in \{x \in A. P\ x\} // R. \text{infinite } X$ 
  from this  $\langle \text{infinite } \{x \in A. P\ x\} \rangle$  show ?thesis
  using  $\langle \bigwedge F. F \in \{x \in A. P\ x\} // R \implies \text{card } F = k \rangle$  card.infinite by auto
qed
qed

```

```

lemma card-equiv-class:
  assumes finite A
  assumes equiv A R
  shows card A = sum card  $(A // R)$ 
proof –
  have  $(\lambda x. \text{True})$  respects R by (simp add: congruentI)
  from  $\langle \text{finite } A \rangle \langle \text{equiv } A\ R \rangle$  this show ?thesis
  using card-equiv-class-restricted[where  $P = \lambda x. \text{True}$ ] by auto
qed

```

```

lemma card-equiv-class-same-size:
  assumes equiv A R
  assumes  $\bigwedge F. F \in A // R \implies \text{card } F = k$ 
  shows card A =  $k * \text{card } (A // R)$ 
proof –
  have  $(\lambda x. \text{True})$  respects R by (simp add: congruentI)
  from  $\langle \text{equiv } A\ R \rangle \langle \bigwedge F. F \in A // R \implies \text{card } F = k \rangle$  this show ?thesis
  using card-equiv-class-restricted-same-size[where  $P = \lambda x. \text{True}$ ] by auto
qed

```

1.3 Additions to FuncSet Theory

```

lemma finite-same-card-bij-on-ext-funcset:
  assumes finite A finite B card A = card B
  shows  $\exists f. f \in A \rightarrow_E B \wedge \text{bij-betw } f\ A\ B$ 
proof –
  from assms obtain f' where  $f': \text{bij-betw } f'\ A\ B$ 
  using finite-same-card-bij by auto
  define f where  $\bigwedge x. f\ x = (\text{if } x \in A \text{ then } f'\ x \text{ else undefined})$ 
  have  $f \in A \rightarrow_E B$ 
  using f' unfolding f-def by (auto simp add: bij-betwE)
  moreover have bij-betw f A B
proof –
  have bij-betw f' A B  $\longleftrightarrow \text{bij-betw } f\ A\ B$ 
  unfolding f-def by (auto intro!: bij-betw-cong)
  from this  $\langle \text{bij-betw } f'\ A\ B \rangle$  show ?thesis by auto
qed
ultimately show ?thesis by auto

```


qed

lemma *card-extensional-funcset*:

assumes *finite A*

shows $\text{card } (A \rightarrow_E B) = \text{card } B \wedge \text{card } A$

using *assms* by (simp add: card-PiE prod-constant)

lemma *bij-betw-implies-inj-on-and-card-eq*:

assumes *finite B*

assumes $f \in A \rightarrow_E B$

shows $\text{bij-betw } f A B \longleftrightarrow \text{inj-on } f A \wedge \text{card } A = \text{card } B$

proof

assume *bij-betw f A B*

from *this* show $\text{inj-on } f A \wedge \text{card } A = \text{card } B$

by (simp add: bij-betw-imp-inj-on bij-betw-same-card)

next

assume $\text{inj-on } f A \wedge \text{card } A = \text{card } B$

from *this* have $\text{inj-on } f A$ and $\text{card } A = \text{card } B$ by *auto*

from $\langle f \in A \rightarrow_E B \rangle$ have $f' A \subseteq B$ by *auto*

from $\langle \text{inj-on } f A \rangle$ have $\text{card } (f' A) = \text{card } A$ by (simp add: card-image)

from $\langle f' A \subseteq B \rangle \langle \text{card } A = \text{card } B \rangle$ *this* have $f' A = B$

by (simp add: $\langle \text{finite } B \rangle$ card-subset-eq)

from $\langle \text{inj-on } f A \rangle$ *this* show *bij-betw f A B* by (rule bij-betw-imageI)

qed

lemma *bij-betw-implies-surj-on-and-card-eq*:

assumes *finite A*

assumes $f \in A \rightarrow_E B$

shows $\text{bij-betw } f A B \longleftrightarrow f' A = B \wedge \text{card } A = \text{card } B$

proof

assume *bij-betw f A B*

show $f' A = B \wedge \text{card } A = \text{card } B$

using $\langle \text{bij-betw } f A B \rangle$ *bij-betw-imp-surj-on bij-betw-same-card* by *blast*

next

assume $f' A = B \wedge \text{card } A = \text{card } B$

from *this* have $f' A = B$ and $\text{card } A = \text{card } B$ by *auto*

from *this* have $\text{inj-on } f A$

by (simp add: $\langle \text{finite } A \rangle$ inj-on-iff-eq-card)

from *this* $\langle f' A = B \rangle$ show *bij-betw f A B* by (rule bij-betw-imageI)

qed

1.4 Additions to Permutations Theory

lemma

assumes $f \in A \rightarrow_E B$ $f' A = B$

assumes *p permutes B* ($\forall x. f' x = p (f x)$)

shows $(\lambda b. \{x \in A. f x = b\})' B = (\lambda b. \{x \in A. f' x = b\})' B$

proof

show $(\lambda b. \{x \in A. f x = b\})' B \subseteq (\lambda b. \{x \in A. f' x = b\})' B$

```

proof
  fix  $X$ 
  assume  $X \in (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B$ 
  from this obtain  $b$  where  $X\text{-eq: } X = \{x \in A. f\ x = b\}$  and  $b \in B$  by blast
  from  $\text{assms}(3, 4)$  have  $\bigwedge x. f\ x = b \longleftrightarrow f'\ x = p\ b$  by (metis permutes-def)
  from  $\langle p \text{ permutes } B \rangle$   $X\text{-eq}$  this have  $X = \{x \in A. f'\ x = p\ b\}$ 
  using Collect-cong by auto
  moreover from  $\langle b \in B \rangle \langle p \text{ permutes } B \rangle$  have  $p\ b \in B$ 
  by (simp add: permutes-in-image)
  ultimately show  $X \in (\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B$  by blast
qed
next
show  $(\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B \subseteq (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B$ 
proof
  fix  $X$ 
  assume  $X \in (\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B$ 
  from this obtain  $b$  where  $X\text{-eq: } X = \{x \in A. f'\ x = b\}$  and  $b \in B$  by blast
  from  $\text{assms}(3, 4)$  have  $\bigwedge x. f'\ x = b \longleftrightarrow f\ x = \text{inv } p\ b$ 
  by (auto simp add: permutes-inverses(1, 2))
  from  $\langle p \text{ permutes } B \rangle$   $X\text{-eq}$  this have  $X = \{x \in A. f\ x = \text{inv } p\ b\}$ 
  using Collect-cong by auto
  moreover from  $\langle b \in B \rangle \langle p \text{ permutes } B \rangle$  have  $\text{inv } p\ b \in B$ 
  by (simp add: permutes-in-image permutes-inv)
  ultimately show  $X \in (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B$  by blast
qed
qed

```

1.5 Additions to List Theory

The theorem *card-lists-length-eq* contains the superfluous assumption *finite A*. Here, we derive that fact without that unnecessary assumption.

lemma *lists-length-eq-Suc-eq-image-Cons*:

$\{xs. \text{set } xs \subseteq A \wedge \text{length } xs = \text{Suc } n\} = (\lambda(x, xs). x\#\text{xs}) \text{ ' } (A \times \{xs. \text{set } xs \subseteq A \wedge \text{length } xs = n\})$
 (*is ?A = ?B*)

proof

show $?A \subseteq ?B$

proof

fix xs

assume $xs \in ?A$

from *this* **show** $xs \in ?B$ **by** (*cases xs*) *auto*

qed

next

show $?B \subseteq ?A$ **by** *auto*

qed

lemma *lists-length-eq-Suc-eq-empty-iff*:

$\{xs. \text{set } xs \subseteq A \wedge \text{length } xs = \text{Suc } n\} = \{\} \longleftrightarrow A = \{\}$

proof (*induct n*)

```

case 0
have {xs. set xs  $\subseteq$  A  $\wedge$  length xs = Suc 0} = {x#[] | x. x  $\in$  A}
proof
  show {[x] | x. x  $\in$  A}  $\subseteq$  {xs. set xs  $\subseteq$  A  $\wedge$  length xs = Suc 0} by auto
next
  show {xs. set xs  $\subseteq$  A  $\wedge$  length xs = Suc 0}  $\subseteq$  {[x] | x. x  $\in$  A}
  proof
    fix xs
    assume xs  $\in$  {xs. set xs  $\subseteq$  A  $\wedge$  length xs = Suc 0}
    from this have set xs  $\subseteq$  A  $\wedge$  length xs = Suc 0 by simp
    from this have  $\exists x. xs = [x] \wedge x \in A$ 
      by (metis Suc-length-conv insert-subset length-0-conv list.set(2))
    from this show xs  $\in$  {[x] | x. x  $\in$  A} by simp
  qed
qed
then show ?case by simp
next
  case (Suc n)
  from this show ?case by (auto simp only: lists-length-eq-Suc-eq-image-Cons)
qed

lemma lists-length-eq-empty-iff:
  {xs. set xs  $\subseteq$  A  $\wedge$  length xs = n} = {}  $\longleftrightarrow$  (A = {}  $\wedge$  n > 0)
proof (cases n)
  case 0
  then show ?thesis by auto
next
  case (Suc n)
  then show ?thesis by (auto simp only: lists-length-eq-Suc-eq-empty-iff)
qed

lemma finite-lists-length-eq-iff:
  finite {xs. set xs  $\subseteq$  A  $\wedge$  length xs = n}  $\longleftrightarrow$  (finite A  $\vee$  n = 0)
proof
  assume finite {xs. set xs  $\subseteq$  A  $\wedge$  length xs = n}
  from this show finite A  $\vee$  n = 0
  proof (induct n)
    case 0
    then show ?case by simp
  next
    case (Suc n)
    have inj ( $\lambda(x, xs). x\#xs$ )
      by (auto intro: inj-onI)
    from this Suc(2) have finite (A  $\times$  {xs. set xs  $\subseteq$  A  $\wedge$  length xs = n})
    using finite-imageD inj-on-subset subset-UNIV lists-length-eq-Suc-eq-image-Cons[of
A n]
      by fastforce
    from this have finite A
      by (cases A = {})
  qed

```

```

      (auto simp only: lists-length-eq-eq-empty-iff dest: finite-cartesian-productD1)
    from this show ?case by auto
  qed
next
  assume finite A  $\vee$   $n = 0$ 
  from this show finite {xs. set xs  $\subseteq$  A  $\wedge$  length xs = n}
    by (auto intro: finite-lists-length-eq)
  qed

lemma card-lists-length-eq:
  shows card {xs. set xs  $\subseteq$  B  $\wedge$  length xs = n} = card B  $\wedge$  n
proof cases
  assume finite B
  then show ?thesis by (rule card-lists-length-eq)
next
  assume infinite B
  then show ?thesis
  proof cases
    assume n = 0
    from this have {xs. set xs  $\subseteq$  B  $\wedge$  length xs = n} = {[]} by auto
    from this  $\langle n = 0 \rangle$  show ?thesis by simp
  next
    assume n  $\neq$  0
    from this  $\langle$ infinite B $\rangle$  have infinite {xs. set xs  $\subseteq$  B  $\wedge$  length xs = n}
      by (simp add: finite-lists-length-eq-iff)
    from this  $\langle$ infinite B $\rangle$  show ?thesis by auto
  qed
qed
qed

```

1.6 Additions to Disjoint Set Theory

```

lemma bij-betw-congI:
  assumes bij-betw f A A'
  assumes  $\forall a \in A. f a = g a$ 
  shows bij-betw g A A'
using assms bij-betw-cong by fastforce

lemma disjoint-family-onI[intro]:
  assumes  $\bigwedge m n. m \in S \implies n \in S \implies m \neq n \implies A m \cap A n = \{\}$ 
  shows disjoint-family-on A S
using assms unfolding disjoint-family-on-def by simp

```

The following lemma is not needed for this development, but is useful and could be moved to Disjoint Set theory or Equiv Relation theory if translated from set partitions to equivalence relations.

```

lemma infinite-partition-on:
  assumes infinite A
  shows infinite {P. partition-on A P}
proof -

```

```

from  $\langle \text{infinite } A \rangle$  obtain  $x$  where  $x \in A$ 
  by (meson finite.intros(1) finite-subset subsetI)
from  $\langle \text{infinite } A \rangle$  have  $\text{infinite } (A - \{x\})$ 
  by (simp add: infinite-remove)
define singletons-except-one
  where singletons-except-one =  $(\lambda a'. (\lambda a. \text{if } a = a' \text{ then } \{a, x\} \text{ else } \{a\})) \text{ ' } (A - \{x\})$ 
have  $\text{infinite } (\text{singletons-except-one ' } (A - \{x\}))$ 
proof –
  have inj-on singletons-except-one  $(A - \{x\})$ 
    unfolding singletons-except-one-def by (rule inj-onI) auto
  from  $\langle \text{infinite } (A - \{x\}) \rangle$  this show ?thesis
    using finite-imageD by blast
qed
moreover have singletons-except-one ' (A - {x})  $\subseteq \{P. \text{partition-on } A \ P\}$ 
proof
  fix  $P$ 
  assume  $P \in \text{singletons-except-one ' } (A - \{x\})$ 
  from this obtain  $a'$  where  $a' \in A - \{x\}$  and  $P: P = \text{singletons-except-one}$ 
 $a'$  by blast
  have partition-on  $A ((\lambda a. \text{if } a = a' \text{ then } \{a, x\} \text{ else } \{a\})) \text{ ' } (A - \{x\})$ 
    using  $\langle x \in A \rangle \langle a' \in A - \{x\} \rangle$  by (auto intro: partition-onI)
  from this have partition-on  $A \ P$ 
    unfolding  $P$  singletons-except-one-def .
  from this show  $P \in \{P. \text{partition-on } A \ P\}$  ..
qed
ultimately show ?thesis by (simp add: infinite-super)
qed

lemma finitely-many-partition-on-iff:
  finite  $\{P. \text{partition-on } A \ P\} \longleftrightarrow \text{finite } A$ 
using finitely-many-partition-on infinite-partition-on by blast

```

1.7 Additions to Multiset Theory

```

lemma mset-set-subseteq-mset-set:
  assumes finite  $B \ A \subseteq B$ 
  shows mset-set  $A \subseteq\# \text{mset-set } B$ 
proof –
  from  $\langle A \subseteq B \rangle \langle \text{finite } B \rangle$  have finite  $A$  using finite-subset by blast
  {
    fix  $x$ 
    have count (mset-set  $A$ )  $x \leq \text{count } (\text{mset-set } B) \ x$ 
      using  $\langle \text{finite } A \rangle \langle \text{finite } B \rangle \langle A \subseteq B \rangle$ 
      by (metis count-mset-set(1, 3) eq-iff subsetCE zero-le-one)
  }
  from this show mset-set  $A \subseteq\# \text{mset-set } B$ 
    using mset-subset-eqI by blast
qed

```

```

lemma mset-set-set-mset:
  assumes  $M \subseteq\# \text{mset-set } A$ 
  shows  $\text{mset-set } (\text{set-mset } M) = M$ 
proof -
  {
    fix  $x$ 
    from  $\langle M \subseteq\# \text{mset-set } A \rangle$  have  $\text{count } M \ x \leq \text{count } (\text{mset-set } A) \ x$ 
      by (simp add: mset-subset-eq-count)
    from this have  $\text{count } (\text{mset-set } (\text{set-mset } M)) \ x = \text{count } M \ x$ 
      by (metis count-eq-zero-iff count-greater-eq-one-iff count-mset-set
        dual-order.antisym dual-order.trans finite-set-mset)
  }
  from this show ?thesis by (simp add: multiset-eq-iff)
qed

lemma mset-set-set-mset':
  assumes  $\forall x. \text{count } M \ x \leq 1$ 
  shows  $\text{mset-set } (\text{set-mset } M) = M$ 
proof -
  {
    fix  $x$ 
    from assms have  $\text{count } M \ x = 0 \vee \text{count } M \ x = 1$  by (auto elim: le-SucE)
    from this have  $\text{count } (\text{mset-set } (\text{set-mset } M)) \ x = \text{count } M \ x$ 
      by (metis count-eq-zero-iff count-mset-set(1,3) finite-set-mset)
  }
  from this show ?thesis by (simp add: multiset-eq-iff)
qed

lemma card-set-mset:
  assumes  $M \subseteq\# \text{mset-set } A$ 
  shows  $\text{card } (\text{set-mset } M) = \text{size } M$ 
using assms
by (metis mset-set-set-mset size-mset-set)

lemma card-set-mset':
  assumes  $\forall x. \text{count } M \ x \leq 1$ 
  shows  $\text{card } (\text{set-mset } M) = \text{size } M$ 
using assms
by (metis mset-set-set-mset' size-mset-set)

lemma count-mset-set-leq:
  assumes finite A
  shows  $\text{count } (\text{mset-set } A) \ x \leq 1$ 
using assms by (metis count-mset-set(1,3) eq-iff zero-le-one)

lemma count-mset-set-leq':
  assumes finite A
  shows  $\text{count } (\text{mset-set } A) \ x \leq \text{Suc } 0$ 

```

using *assms count-mset-set-leq* **by** *fastforce*

lemma *msubset-mset-set-iff*:

assumes *finite A*

shows $\text{set-mset } M \subseteq A \wedge (\forall x. \text{count } M \ x \leq 1) \longleftrightarrow (M \subseteq\# \text{mset-set } A)$

proof

assume $\text{set-mset } M \subseteq A \wedge (\forall x. \text{count } M \ x \leq 1)$

from *this assms* **show** $M \subseteq\# \text{mset-set } A$

by (*metis count-inI count-mset-set(1) le0 mset-subset-eqI subsetCE*)

next

assume $M \subseteq\# \text{mset-set } A$

from *this assms* **have** $\text{set-mset } M \subseteq A$

using *mset-subset-eqD* **by** *fastforce*

moreover {

fix *x*

from $\langle M \subseteq\# \text{mset-set } A \rangle$ **have** $\text{count } M \ x \leq \text{count } (\text{mset-set } A) \ x$

by (*simp add: mset-subset-eq-count*)

from *this* $\langle \text{finite } A \rangle$ **have** $\text{count } M \ x \leq 1$

by (*meson count-mset-set-leq le-trans*)

}

ultimately show $\text{set-mset } M \subseteq A \wedge (\forall x. \text{count } M \ x \leq 1)$ **by** *simp*

qed

lemma *image-mset-fun-upd*:

assumes $x \notin\# M$

shows $\text{image-mset } (f(x := y)) \ M = \text{image-mset } f \ M$

using *assms* **by** (*induct M*) *auto*

1.8 Additions to Number Partitions Theory

lemma *Partition-diag*:

shows $\text{Partition } n \ n = 1$

by (*cases n*) (*auto simp only: Partition-diag Partition.simps(1)*)

1.9 Cardinality Theorems with Iverson Function

definition *iverson* :: $\text{bool} \Rightarrow \text{nat}$

where

$\text{iverson } b = (\text{if } b \text{ then } 1 \text{ else } 0)$

lemma *card-partition-on-size1-eq-iverson*:

assumes *finite A*

shows $\text{card } \{P. \text{partition-on } A \ P \wedge \text{card } P \leq k \wedge (\forall X \in P. \text{card } X = 1)\} = \text{iverson } (\text{card } A \leq k)$

proof (*cases card A ≤ k*)

case *True*

from *this* $\langle \text{finite } A \rangle$ **show** *?thesis*

unfolding *iverson-def*

using *card-partition-on-size1-eq-1* **by** *fastforce*

next

```

    case False
    from this <finite A> show ?thesis
      unfolding iverson-def
      using card-partition-on-size1-eq-0 by fastforce
qed

lemma card-number-partitions-with-only-parts-1:
  card {N. (∀ n. n ∈ # N ⟶ n = 1) ∧ number-partition n N ∧ size N ≤ x} =
  iverson (n ≤ x)
proof -
  show ?thesis
  proof cases
    assume n ≤ x
    from this show ?thesis
      using card-number-partitions-with-only-parts-1-eq-1
      unfolding iverson-def by auto
  next
    assume ¬ n ≤ x
    from this show ?thesis
      using card-number-partitions-with-only-parts-1-eq-0
      unfolding iverson-def by auto
  qed
qed
end

```

2 Main Observations on Operations and Permutations

```

theory Twelvefold-Way-Core
imports Preliminaries
begin

```

2.1 Range Multiset

2.1.1 Existence of a Suitable Finite Function

```

lemma obtain-function:
  assumes finite A
  assumes size M = card A
  shows ∃ f. image-mset f (mset-set A) = M
using assms
proof (induct arbitrary: M rule: finite-induct)
  case empty
  from this show ?case by simp
next
  case (insert x A)
  from insert(1,2,4) have size M > 0
  by (simp add: card-gt-0-iff)

```



```

from this obtain  $y$  where  $y \in \# M$ 
  using gr0-implies-Suc size-eq-Suc-imp-elem by blast
from insert(1,2,4) this have  $\text{size } (M - \{\#y\# \}) = \text{card } A$ 
  by (simp add: Diff-insert-absorb card-Diff-singleton-if insertI1 size-Diff-submset)
from insert.hyps this obtain  $f'$  where  $\text{image-mset } f' (\text{mset-set } A) = M - \{\#y\# \}$  by blast
from this have  $\text{image-mset } (f'(x := y)) (\text{mset-set } (\text{insert } x A)) = M$ 
  using  $\langle \text{finite } A \rangle \langle x \notin A \rangle \langle y \in \# M \rangle$  by (simp add: image-mset-fun-upd)
from this show ?case by blast
qed

```

lemma *obtain-function-on-ext-funcset:*

```

assumes finite A
assumes  $\text{size } M = \text{card } A$ 
shows  $\exists f \in A \rightarrow_E \text{set-mset } M. \text{image-mset } f (\text{mset-set } A) = M$ 
proof –
  obtain  $f$  where range-eq-M:  $\text{image-mset } f (\text{mset-set } A) = M$ 
    using obtain-function  $\langle \text{finite } A \rangle \langle \text{size } M = \text{card } A \rangle$  by blast
  let  $?f = \lambda x. \text{if } x \in A \text{ then } f \text{ else undefined}$ 
  have  $?f \in A \rightarrow_E \text{set-mset } M$ 
    using range-eq-M  $\langle \text{finite } A \rangle$  by auto
  moreover have  $\text{image-mset } ?f (\text{mset-set } A) = M$ 
    using range-eq-M  $\langle \text{finite } A \rangle$  by (auto intro: multiset.map-cong0)
  ultimately show ?thesis by auto
qed

```

2.1.2 Existence of Permutation

lemma *image-mset-eq-implies-bij-betw:*

```

fixes  $f :: 'a1 \Rightarrow 'b$  and  $f' :: 'a2 \Rightarrow 'b$ 
assumes finite A finite A'
assumes mset-eq:  $\text{image-mset } f (\text{mset-set } A) = \text{image-mset } f' (\text{mset-set } A')$ 
obtains bij where bij-betw bij A A' and  $\forall x \in A. f x = f' (\text{bij } x)$ 
proof –
  from  $\langle \text{finite } A \rangle$  have [simp]:  $\text{finite } \{a \in A. f a = (b::'b)\}$  for  $b$  by auto
  from  $\langle \text{finite } A' \rangle$  have [simp]:  $\text{finite } \{a \in A'. f' a = (b::'b)\}$  for  $b$  by auto
  have  $f' A = f' A'$ 
  proof –
    have  $f' A = f' (\text{set-mset } (\text{mset-set } A))$  using  $\langle \text{finite } A \rangle$  by simp
    also have  $\dots = f' (\text{set-mset } (\text{mset-set } A'))$ 
      by (metis mset-eq multiset.set-map)
    also have  $\dots = f' A'$  using  $\langle \text{finite } A' \rangle$  by simp
    finally show ?thesis .
  qed

```

have $\forall b \in (f' A). \exists \text{bij}. \text{bij-betw } \text{bij } \{a \in A. f a = b\} \{a \in A'. f' a = b\}$

proof

fix b

from *mset-eq* **have**

$\text{count } (\text{image-mset } f (\text{mset-set } A)) b = \text{count } (\text{image-mset } f' (\text{mset-set } A')) b$

```

by simp
  from this have card {a ∈ A. f a = b} = card {a ∈ A'. f' a = b}
  using ⟨finite A⟩ ⟨finite A'⟩
  by (simp add: count-image-mset-eq-card-vimage)
  from this show ∃ bij. bij-betw bij {a ∈ A. f a = b} {a ∈ A'. f' a = b}
  by (intro finite-same-card-bij) simp-all
qed
from bchoice [OF this]
obtain bij where bij: ∀ b ∈ f ' A. bij-betw (bij b) {a ∈ A. f a = b} {a ∈ A'. f' a
= b}
  by auto
define bij' where bij' = (λa. bij (f a) a)
have bij-betw bij' A A'
proof -
  have disjoint-family-on (λi. {a ∈ A'. f' a = i}) (f ' A)
  unfolding disjoint-family-on-def by auto
  moreover have bij-betw (λa. bij (f a) a) {a ∈ A. f a = b} {a ∈ A'. f' a = b}
if b: b ∈ f ' A for b
  using bij b by (subst bij-betw-cong[where g=bij b]) auto
  ultimately have bij-betw (λa. bij (f a) a) (⋃ b ∈ f ' A. {a ∈ A. f a = b}) (⋃ b ∈ f
' A. {a ∈ A'. f' a = b})
  by (rule bij-betw-UNION-disjoint)
  moreover have (⋃ b ∈ f ' A. {a ∈ A. f a = b}) = A by auto
  moreover have (⋃ b ∈ f ' A. {a ∈ A'. f' a = b}) = A' using ⟨f ' A = f' ' A'⟩
by auto
  ultimately show bij-betw bij' A A'
  unfolding bij'-def by (subst bij-betw-cong[where g=(λa. bij (f a) a)]) auto
qed
moreover from bij have ∀ x ∈ A. f x = f' (bij' x)
  unfolding bij'-def using bij-betwE by fastforce
ultimately show ?thesis by (rule that)
qed

lemma image-mset-eq-implies-permutes:
  fixes f :: 'a ⇒ 'b
  assumes finite A
  assumes mset-eq: image-mset f (mset-set A) = image-mset f' (mset-set A)
  obtains p where p permutes A and ∀ x ∈ A. f x = f' (p x)
proof -
  from assms obtain b where bij-betw b A A and ∀ x ∈ A. f x = f' (b x)
  using image-mset-eq-implies-bij-betw by blast
  define p where p = (λa. if a ∈ A then b a else a)
  have p permutes A
  proof (rule bij-imp-permutes)
    show bij-betw p A A
    unfolding p-def by (simp add: ⟨bij-betw b A A⟩ bij-betw-cong)
  next
    fix x
    assume x ∉ A

```

from *this* show $p\ x = x$
 unfolding $p\text{-def}$ by *simp*
 qed
 moreover from $\langle \forall x \in A. f\ x = f'\ (b\ x) \rangle$ have $\forall x \in A. f\ x = f'\ (p\ x)$
 unfolding $p\text{-def}$ by *simp*
 ultimately show *?thesis* by (rule *that*)
 qed

2.2 Domain Partition

2.2.1 Existence of a Suitable Finite Function

lemma *obtain-function-with-partition*:

assumes *finite A finite B*
 assumes *partition-on A P*
 assumes *card P ≤ card B*
 shows $\exists f \in A \rightarrow_E B. (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\} = P$
 proof –
 obtain g' where *bij-betw $g' P (g' \text{ ' } P)$ and $g' \text{ ' } P \subseteq B$*
 by (*meson assms card-le-inj finite-elements inj-on-imp-bij-betw*)
 define f where $\bigwedge a. f\ a = (\text{if } a \in A \text{ then } g'\ (THE\ X. a \in X \wedge X \in P) \text{ else } undefined)$
 have $f \in A \rightarrow_E B$
 unfolding $f\text{-def}$
 using $\langle g' \text{ ' } P \subseteq B \rangle$ *assms(3) partition-on-the-part-mem* by *fastforce*
 moreover have $(\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\} = P$
 proof
 show $(\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\} \subseteq P$
 proof
 fix X
 assume $X: X \in (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}$
 from *this* obtain b where $b \in B$ and $X = \{x' \in A. f\ x' = b\}$ by *auto*
 from *this* X obtain a where $a \in A$ and $a \in X$ and $f\ a = b$ by *blast*
 have $(THE\ X. a \in X \wedge X \in P) \in P$
 using $\langle a \in A \rangle \langle \text{partition-on } A\ P \rangle$ by (*simp add: partition-on-the-part-mem*)
 from $\langle X = \{x' \in A. f\ x' = b\} \rangle$ have $X\text{-eq1}: X = \{x' \in A. g'\ (THE\ X. x' \in X \wedge X \in P) = b\}$
 unfolding $f\text{-def}$ by *auto*
 also have $\dots = \{x' \in A. (THE\ X. x' \in X \wedge X \in P) = inv\text{-into } P\ g'\ b\}$
 proof –
 {
 fix x'
 assume $x' \in A$
 have $(THE\ X. x' \in X \wedge X \in P) \in P$
 using $\langle \text{partition-on } A\ P \rangle \langle x' \in A \rangle$ by (*simp add: partition-on-the-part-mem*)
 from $X\text{-eq1}$ $\langle a \in X \rangle$ have $g'\ (THE\ X. a \in X \wedge X \in P) = b$
 unfolding $f\text{-def}$ by *auto*
 from *this* $\langle (THE\ X. a \in X \wedge X \in P) \in P \rangle$ have $b \in g' \text{ ' } P$ by *auto*
 have $(g'\ (THE\ X. x' \in X \wedge X \in P) = b) \longleftrightarrow ((THE\ X. x' \in X \wedge X \in P) = inv\text{-into } P\ g'\ b)$

```

proof –
  from  $\langle (THE\ X.\ x' \in X \wedge X \in P) \in P \rangle$ 
  have  $(g' (THE\ X.\ x' \in X \wedge X \in P) = b) \longleftrightarrow (inv\text{-}into\ P\ g' (g' (THE\ X.\ x' \in X \wedge X \in P)) = inv\text{-}into\ P\ g' b)$ 
  using  $\langle b \in g' \text{' } P \rangle$  by (auto intro: inv-into-injective)
  moreover have  $inv\text{-}into\ P\ g' (g' (THE\ X.\ x' \in X \wedge X \in P)) = (THE\ X.\ x' \in X \wedge X \in P)$ 
  using  $\langle bij\text{-}betw\ g' P\ (g' \text{' } P) \rangle$   $\langle (THE\ X.\ x' \in X \wedge X \in P) \in P \rangle$ 
  by (simp add: bij-betw-inv-into-left)
  ultimately show ?thesis by simp
qed
}
from this show ?thesis by auto
qed
finally have  $X\text{-}eq: X = \{x' \in A.\ (THE\ X.\ x' \in X \wedge X \in P) = inv\text{-}into\ P\ g' b\}$ .
moreover have  $inv\text{-}into\ P\ g' b \in P$ 
proof –
  from  $X\text{-}eq$  have  $eq: inv\text{-}into\ P\ g' b = (THE\ X.\ a \in X \wedge X \in P)$ 
  using  $\langle a \in X \rangle$   $\langle a \in A \rangle$  by auto
  from this show ?thesis
  using  $\langle (THE\ X.\ a \in X \wedge X \in P) \in P \rangle$  by simp
qed
ultimately have  $X = inv\text{-}into\ P\ g' b$ 
  using partition-on-all-in-part-eq-part[OF  $\langle partition\text{-}on\ A\ P \rangle$ ] by blast
from this  $\langle inv\text{-}into\ P\ g' b \in P \rangle$  show  $X \in P$  by blast
qed
next
show  $P \subseteq (\lambda b.\ \{x \in A.\ f\ x = b\}) \text{' } B - \{\{\}\}$ 
proof
  fix  $X$ 
  assume  $X \in P$ 
  from  $assms(3)$  this have  $X \neq \{\}$ 
  by (auto elim: partition-onE)
  moreover have  $X \in (\lambda b.\ \{x \in A.\ f\ x = b\}) \text{' } B$ 
proof
  show  $g' X \in B$ 
  using  $\langle X \in P \rangle$   $\langle g' \text{' } P \subseteq B \rangle$  by blast
  show  $X = \{x \in A.\ f\ x = g' X\}$ 
proof
  show  $X \subseteq \{x \in A.\ f\ x = g' X\}$ 
  proof
    fix  $x$ 
    assume  $x \in X$ 
    from this have  $x \in A$ 
    using  $\langle X \in P \rangle$   $assms(3)$  by (fastforce elim: partition-onE)
    have  $(THE\ X.\ x \in X \wedge X \in P) = X$ 
    using  $\langle X \in P \rangle$   $\langle x \in X \rangle$   $assms(3)$  partition-on-the-part-eq by fastforce
    from this  $\langle x \in A \rangle$  have  $f\ x = g' X$ 
  
```

```

      unfolding f-def by auto
      from this  $\langle x \in A \rangle$  show  $x \in \{x \in A. f\ x = g'\ X\}$  by auto
    qed
  next
  show  $\{x \in A. f\ x = g'\ X\} \subseteq X$ 
  proof
    fix x
    assume  $x \in \{x \in A. f\ x = g'\ X\}$ 
    from this have  $x \in A$  and g-eq:  $g'\ (THE\ X. x \in X \wedge X \in P) = g'\ X$ 
      unfolding f-def by auto
    from  $\langle x \in A \rangle$  have  $(THE\ X. x \in X \wedge X \in P) \in P$ 
      using assms(3) by (simp add: partition-on-the-part-mem)
    from this g-eq have  $(THE\ X. x \in X \wedge X \in P) = X$ 
      using  $\langle X \in P \rangle$  <bij-betw  $g'\ P\ (g'\ ' P)$ >
      by (metis bij-betw-inv-into-left)
    from this  $\langle x \in A \rangle$  assms(3) show  $x \in X$ 
      using partition-on-in-the-unique-part by fastforce
    qed
  qed
  qed
  ultimately show  $X \in (\lambda b. \{x \in A. f\ x = b\})\ ' B - \{\{\}\}$ 
    by auto
  qed
  qed
  ultimately show ?thesis by blast
qed

```

2.2.2 Equality under Permutation Application

lemma *permutes-implies-inv-image-on-eq*:

```

  assumes p permutes B
  shows  $(\lambda b. \{x \in A. p\ (f\ x) = b\})\ ' B = (\lambda b. \{x \in A. f\ x = b\})\ ' B$ 
  proof -
    have  $\forall b \in B. \forall x \in A. p\ (f\ x) = b \longleftrightarrow f\ x = inv\ p\ b$ 
      using  $\langle p\ permutes\ B \rangle$  by (auto simp add: permutes-inverses)
    from this have  $(\lambda b. \{x \in A. p\ (f\ x) = b\})\ ' B = (\lambda b. \{x \in A. f\ x = inv\ p\ b\})\ ' B$ 
      using image-cong by blast
    also have  $\dots = (\lambda b. \{x \in A. f\ x = b\})\ ' inv\ p\ ' B$ 
      by (auto simp add: image-comp)
    also have  $\dots = (\lambda b. \{x \in A. f\ x = b\})\ ' B$ 
      by (simp add:  $\langle p\ permutes\ B \rangle$  permutes-inv permutes-image)
    finally show ?thesis .
  qed

```

2.2.3 Existence of Permutation

lemma *the-elem*:

```

  assumes  $f \in A \rightarrow_E B$   $f' \in A \rightarrow_E B$ 

```

assumes *partitions-eq*: $(\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\} = (\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B - \{\{\}\}$
assumes $x \in A$
shows *the-elem* $(f \text{ ' } \{x \in A. f'\ x = f'\ x\}) = f\ x$
proof –
from $\langle x \in A \rangle$ **have** $x: x \in \{x' \in A. f'\ x' = f'\ x\}$ **by** *blast*
have $f'\ x \in B$
using $\langle x \in A \rangle \langle f' \in A \rightarrow_E B \rangle$ **by** *blast*
from *this* **have** $\{x' \in A. f'\ x' = f'\ x\} \in (\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B - \{\{\}\}$
using $\langle x \in A \rangle$ **by** *blast*
from *this* **have** $\{x' \in A. f'\ x' = f'\ x\} \in (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}$
using *partitions-eq* **by** *blast*
from *this* **obtain** b **where** *eq*: $\{x' \in A. f'\ x' = f'\ x\} = \{x' \in A. f\ x' = b\}$ **by** *blast*
also from x **this show** *the-elem* $(f \text{ ' } \{x' \in A. f'\ x' = f'\ x\}) = f\ x$
by (*metis* (*mono-tags*, *lifting*) *empty-iff mem-Collect-eq the-elem-image-unique*)
qed

lemma *the-elem-eq*:
assumes $f \in A \rightarrow_E B$
assumes $b \in f \text{ ' } A$
shows *the-elem* $(f \text{ ' } \{x' \in A. f\ x' = b\}) = b$
proof –
from $\langle b \in f \text{ ' } A \rangle$ **obtain** $a \in A$ **and** $b = f\ a$ **by** *blast*
from *this* **show** *the-elem* $(f \text{ ' } \{x' \in A. f\ x' = b\}) = b$
using *the-elem*[*OF* $\langle f \in A \rightarrow_E B \rangle \langle f \in A \rightarrow_E B \rangle$] **by** *simp*
qed

lemma *partitions-eq-implies*:
assumes $f \in A \rightarrow_E B$ $f' \in A \rightarrow_E B$
assumes *partitions-eq*: $(\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\} = (\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B - \{\{\}\}$
assumes $x \in A$ $x' \in A$
assumes $f\ x = f'\ x'$
shows $f'\ x = f'\ x'$
proof –
have $f\ x \in B$ **and** $x \in \{a \in A. f\ a = f\ x\}$ **and** $x' \in \{a \in A. f'\ a = f\ x\}$
using $\langle f \in A \rightarrow_E B \rangle \langle x \in A \rangle \langle x' \in A \rangle \langle f\ x = f'\ x' \rangle$ **by** *auto*
moreover have $\{a \in A. f'\ a = f\ x\} \in (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}$
using $\langle f\ x \in B \rangle \langle x \in \{a \in A. f'\ a = f\ x\} \rangle$ **by** *auto*
ultimately obtain b **where** $x \in \{a \in A. f'\ a = b\}$ **and** $x' \in \{a \in A. f'\ a = b\}$
using *partitions-eq* **by** (*metis* (*no-types*, *lifting*) *Diff-iff imageE*)
from *this* **show** $f'\ x = f'\ x'$ **by** *auto*
qed

lemma *card-domain-partitions*:
assumes $f \in A \rightarrow_E B$
assumes *finite* B
shows *card* $((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}) = \text{card}\ (f \text{ ' } A)$

```

proof –
  note [simp] = the-elem-eq[OF ⟨f ∈ A →E B⟩]
  have bij-betw (λX. the-elem (f ‘ X)) ((λb. {x ∈ A. f x = b}) ‘ B – {{})) (f ‘ A)
  proof (rule bij-betw-imageI)
    show inj-on (λX. the-elem (f ‘ X)) ((λb. {x ∈ A. f x = b}) ‘ B – {{}))
    proof (rule inj-onI)
      fix X X'
      assume X: X ∈ (λb. {x ∈ A. f x = b}) ‘ B – {{}
      assume X': X' ∈ (λb. {x ∈ A. f x = b}) ‘ B – {{}
      assume eq: the-elem (f ‘ X) = the-elem (f ‘ X')
      from X obtain b where b ∈ B and X-eq: X = {x ∈ A. f x = b} by blast
      from X this have b ∈ f ‘ A
      using Collect-empty-eq Diff-iff image-iff insertCI by auto
      from X' obtain b' where b' ∈ B and X'-eq: X' = {x ∈ A. f x = b'} by
blast
      from X' this have b' ∈ f ‘ A
      using Collect-empty-eq Diff-iff image-iff insertCI by auto
      from X-eq X'-eq eq ⟨λb. b ∈ f ‘ A ⇒ the-elem (f ‘ {x' ∈ A. f x' = b}) = b⟩
⟨b ∈ f ‘ A⟩ ⟨b' ∈ f ‘ A⟩
      have b = b' by auto
      from this show X = X'
      using X-eq X'-eq by simp
    qed
  show (λX. the-elem (f ‘ X)) ‘ ((λb. {x ∈ A. f x = b}) ‘ B – {{})) = f ‘ A
  proof
    show (λX. the-elem (f ‘ X)) ‘ ((λb. {x ∈ A. f x = b}) ‘ B – {{})) ⊆ f ‘ A
    using ⟨λb. b ∈ f ‘ A ⇒ the-elem (f ‘ {x' ∈ A. f x' = b}) = b⟩ by auto
  next
    show f ‘ A ⊆ (λX. the-elem (f ‘ X)) ‘ ((λb. {x ∈ A. f x = b}) ‘ B – {{}))
    proof
      fix b
      assume b ∈ f ‘ A
      from this have b = the-elem (f ‘ {x ∈ A. f x = b})
      using ⟨λb. b ∈ f ‘ A ⇒ the-elem (f ‘ {x' ∈ A. f x' = b}) = b⟩ by auto
      moreover from ⟨b ∈ f ‘ A⟩ have {x ∈ A. f x = b} ∈ (λb. {x ∈ A. f x =
b}) ‘ B – {{}
      using ⟨f ∈ A →E B⟩ by auto
      ultimately show b ∈ (λX. the-elem (f ‘ X)) ‘ ((λb. {x ∈ A. f x = b}) ‘ B
– {{})) ..
    qed
  qed
qed
from this show ?thesis by (rule bij-betw-same-card)
qed

```

```

lemma partitions-eq-implies-permutes:
  assumes f ∈ A →E B f' ∈ A →E B
  assumes finite B
  assumes partitions-eq: (λb. {x ∈ A. f x = b}) ‘ B – {{} = (λb. {x ∈ A. f' x

```

```

= b}) ' B - { {} }
shows  $\exists p. p \text{ permutes } B \wedge (\forall x \in A. f x = p (f' x))$ 
proof -
  have card-eq: card (f' ' A) = card (f ' A)
    using card-domain-partitions[OF  $\langle f \in A \rightarrow_E B \rangle \langle \text{finite } B \rangle$ ]
    using card-domain-partitions[OF  $\langle f' \in A \rightarrow_E B \rangle \langle \text{finite } B \rangle$ ]
    using partitions-eq by simp
  have f' ' A  $\subseteq$  B f' ' A  $\subseteq$  B
    using  $\langle f \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle$  by auto
  from this card-eq have card (B - f' ' A) = card (B - f ' A)
    using  $\langle \text{finite } B \rangle$  by (auto simp add: card-Diff-subset finite-subset)
  from this obtain p' where bij-betw p' (B - f' ' A) (B - f ' A)
    using  $\langle \text{finite } B \rangle$  by (metis finite-same-card-bij finite-Diff)
  from this have p' ' (B - f' ' A) = (B - f ' A)
    by (simp add: bij-betw-imp-surj-on)
  define p where  $\bigwedge b. p b = (\text{if } b \in B \text{ then } (\text{if } b \in f' ' A \text{ then the-elem } (f' ' \{x \in A. f' x = b\}) \text{ else } p' b) \text{ else } b)$ 
  have  $\forall x \in A. f x = p (f' x)$ 
  proof
    fix x
    assume x  $\in$  A
    from this partitions-eq have the-elem (f' ' {xa  $\in$  A. f' xa = f' x}) = f x
      using the-elem[OF  $\langle f \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle$ ] by auto
    from this show f x = p (f' x)
      using  $\langle x \in A \rangle$  p-def  $\langle f' \in A \rightarrow_E B \rangle$  by auto
  qed
  moreover have p permutes B
  proof (rule bij-imp-permutes)
    let ?invp =  $\lambda b. \text{if } b \in f' ' A \text{ then the-elem } (f' ' \{x \in A. f x = b\}) \text{ else } b$ 
    note [simp] = the-elem[OF  $\langle f \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle$  partitions-eq]
    show bij-betw p B B
    proof (rule bij-betw-imageI)
      show p ' B = B
    proof
      have ( $\lambda b. \text{the-elem } (f' ' \{x \in A. f' x = b\})$ ) ' (f' ' A)  $\subseteq$  B
        using  $\langle f \in A \rightarrow_E B \rangle$  by auto
      from  $\langle p' ' (B - f' ' A) = (B - f ' A) \rangle$  this show p ' B  $\subseteq$  B
        unfolding p-def  $\langle f \in A \rightarrow_E B \rangle$  by force
    next
      show B  $\subseteq$  p ' B
    proof
      fix b
      assume b  $\in$  B
      show b  $\in$  p ' B
    proof (cases b  $\in$  f' ' A)
      assume b  $\notin$  f' ' A
      note  $\langle p' ' (B - f' ' A) = (B - f ' A) \rangle$ 
      from this  $\langle b \in B \rangle \langle b \notin f' ' A \rangle$  show ?thesis
        unfolding p-def by auto
    qed
  qed
  qed

```



```

next
  assume  $b \in f' \text{ ' } A$ 
  from this  $\langle \forall x \in A. f \ x = p \ (f' \ x) \rangle \langle b \in B \rangle$  show ?thesis
  using  $\langle f' \in A \rightarrow_E B \rangle$  by auto
qed
qed
qed
next
show inj-on  $p \ B$ 
proof (rule inj-onI)
  fix  $b \ b'$ 
  assume  $b \in B \ b' \in B \ p \ b = p \ b'$ 
  have  $b \in f' \text{ ' } A \longleftrightarrow b' \in f' \text{ ' } A$ 
  proof –
    have  $b \in f' \text{ ' } A \longleftrightarrow p \ b \in f' \text{ ' } A$ 
    unfolding p-def using  $\langle b \in B \rangle \langle p' \text{ ' } (B - f' \text{ ' } A) = B - f' \text{ ' } A \rangle$  by auto
    also have  $p \ b \in f' \text{ ' } A \longleftrightarrow p \ b' \in f' \text{ ' } A$ 
    using  $\langle p \ b = p \ b' \rangle$  by simp
    also have  $p \ b' \in f' \text{ ' } A \longleftrightarrow b' \in f' \text{ ' } A$ 
    unfolding p-def using  $\langle b' \in B \rangle \langle p' \text{ ' } (B - f' \text{ ' } A) = B - f' \text{ ' } A \rangle$  by auto
    finally show ?thesis .
  qed
  from this have  $(b \in f' \text{ ' } A \wedge b' \in f' \text{ ' } A) \vee (b \notin f' \text{ ' } A \wedge b' \notin f' \text{ ' } A)$  by
blast
  from this show  $b = b'$ 
  proof
    assume  $b \in f' \text{ ' } A \wedge b' \in f' \text{ ' } A$ 
    from this obtain  $a \ a'$  where  $a \in A \ b = f' \ a$  and  $a' \in A \ b' = f' \ a'$  by
auto
    from this  $\langle b \in B \rangle \langle b' \in B \rangle$  have  $p \ b = f \ a \ p \ b' = f \ a'$ 
    unfolding p-def by auto
    from this  $\langle p \ b = p \ b' \rangle$  have  $f \ a = f \ a'$  by simp
    from this have  $f' \ a = f' \ a'$ 
    using partitions-eq-implies[OF  $\langle f \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle$  partitions-eq]
    using  $\langle a \in A \rangle \langle a' \in A \rangle$  by blast
    from this show  $b = b'$ 
    using  $\langle b' = f' \ a' \rangle \langle b = f' \ a \rangle$  by simp
  next
    assume  $b \notin f' \text{ ' } A \wedge b' \notin f' \text{ ' } A$ 
    from this  $\langle b \in B \rangle \langle b' \in B \rangle$  have  $p \ b' = p' \ b' \ p \ b = p' \ b$ 
    unfolding p-def by auto
    from this  $\langle p \ b = p \ b' \rangle$  have  $p' \ b = p' \ b'$  by simp
    moreover have  $b \in B - f' \text{ ' } A \ b' \in B - f' \text{ ' } A$ 
    using  $\langle b \in B \rangle \langle b' \in B \rangle \langle b \notin f' \text{ ' } A \wedge b' \notin f' \text{ ' } A \rangle$  by auto
    ultimately show  $b = b'$ 
    using bij-betw  $p' \ - \rightarrow$  by (metis bij-betw-inv-into-left)
  qed
qed
qed
qed

```

```

next
  fix x
  assume  $x \notin B$ 
  from this show  $p\ x = x$ 
    using  $\langle f' \in A \rightarrow_E B \rangle$  p-def by auto
  qed
  ultimately show ?thesis by blast
qed

```

2.3 Number Partition of Range

2.3.1 Existence of a Suitable Finite Function

```

lemma obtain-partition:
  assumes finite A
  assumes number-partition (card A) N
  shows  $\exists P. \text{partition-on } A\ P \wedge \text{image-mset card (mset-set } P) = N$ 
using assms
proof (induct N arbitrary: A)
  case empty
  from this have  $A = \{\}$ 
    unfolding number-partition-def by auto
  from this have partition-on  $A\ \{\}$  by (simp add: partition-on-empty)
  moreover have image-mset card (mset-set  $\{\}$ ) =  $\{\#\}$  by simp
  ultimately show ?case by blast
next
  case (add x N)
  from add.prems(2) have  $0 \notin \# \text{ add-mset } x\ N$  and sum-mset (add-mset  $x\ N) =$ 
card A
    unfolding number-partition-def by auto
  from this have  $x \leq \text{card } A$  by auto
  from this obtain X where  $X \subseteq A$  and card  $X = x$ 
    using subset-with-given-card-exists by auto
  from this have  $X \neq \{\}$ 
    using  $\langle 0 \notin \# \text{ add-mset } x\ N \rangle \langle \text{finite } A \rangle$  by auto
  have sum-mset  $N = \text{card } (A - X)$ 
    using  $\langle \text{sum-mset (add-mset } x\ N) = \text{card } A \rangle \langle \text{card } X = x \rangle \langle X \subseteq A \rangle$ 
    by (metis add.commute add.prems(1) add-diff-cancel-right' card-Diff-subset
infinite-super sum-mset.add-mset)
  from this  $\langle 0 \notin \# \text{ add-mset } x\ N \rangle$  have number-partition (card (A - X)) N
    unfolding number-partition-def by auto
  from this obtain P where partition-on  $(A - X)\ P$  and eq-N: image-mset card
(mset-set P) = N
    using add.hyps  $\langle \text{finite } A \rangle$  by auto
  from  $\langle \text{partition-on } (A - X)\ P \rangle$  have finite P
    using  $\langle \text{finite } A \rangle$  finite-elements by blast
  from  $\langle \text{partition-on } (A - X)\ P \rangle$  have  $X \notin P$ 
    using  $\langle X \neq \{\} \rangle$  partition-onD1 by fastforce
  have partition-on  $A\ (\text{insert } X\ P)$ 
    using  $\langle \text{partition-on } (A - X)\ P \rangle \langle X \subseteq A \rangle \langle X \neq \{\} \rangle$ 

```

by (rule partition-on-insert')
 moreover have image-mset card (mset-set (insert X P)) = add-mset x N
 using eq-N ⟨card X = x⟩ ⟨finite P⟩ ⟨X ∉ P⟩ by simp
 ultimately show ?case by blast
 qed

lemma obtain-extensional-function-from-number-partition:
 assumes finite A finite B
 assumes number-partition (card A) N
 assumes size N ≤ card B
 shows $\exists f \in A \rightarrow_E B. \text{image-mset } (\lambda X. \text{card } X) (\text{mset-set } ((\lambda b. \{x \in A. f\ x = b\}))) \text{ ' } B - \{\{\}\} = N$
proof –
 obtain P where partition-on A P and eq-N: image-mset card (mset-set P) = N
 using assms obtain-partition by blast
 from eq-N[symmetric] ⟨size N ≤ card B⟩ have card P ≤ card B by simp
 from ⟨partition-on A P⟩ this obtain f where $f \in A \rightarrow_E B$
 and eq-P: $(\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\} = P$
 using obtain-function-with-partition[OF ⟨finite A⟩ ⟨finite B⟩] by blast
 have image-mset (λX. card X) (mset-set (((λb. {x ∈ A. f x = b})) ' B - {{})))
 = N
 using eq-P eq-N by simp
 from this ⟨f ∈ A →_E B⟩ show ?thesis by auto
 qed

2.3.2 Equality under Permutation Application

lemma permutes-implies-multiset-of-partition-cards-eq:
 assumes p_A permutes A p_B permutes B
 shows image-mset card (mset-set ((λb. {x ∈ A. p_B (f' (p_A x)) = b}) ' B - {{}))) = image-mset card (mset-set ((λb. {x ∈ A. f' x = b}) ' B - {{})))
proof –
 have inj-on ((\cdot) (inv p_A)) ((λb. {x ∈ A. f' x = b}) ' B - {{}))
 by (meson ⟨p_A permutes A⟩ inj-image-eq-iff inj-onI permutes-surj surj-imp-inj-inv)
 have image-mset card (mset-set ((λb. {x ∈ A. p_B (f' (p_A x)) = b}) ' B - {{})))
 =
 image-mset card (mset-set ((λX. inv p_A ' X) ' ((λb. {x ∈ A. f' x = b}) ' B - {{}))))
proof –
 have (λb. {x ∈ A. p_B (f' (p_A x)) = b}) ' B - {{} = (λb. {x ∈ A. f' (p_A x) = b}) ' B - {{}
 = b}) ' B - {{}
 using permutes-implies-inv-image-on-eq[OF ⟨p_B permutes B⟩] by metis
 also have ... = (λb. inv p_A ' {x ∈ A. f' x = b}) ' B - {{}
proof –
 have {x ∈ A. f' (p_A x) = b} = inv p_A ' {x ∈ A. f' x = b} for b
proof
 show {x ∈ A. f' (p_A x) = b} ⊆ inv p_A ' {x ∈ A. f' x = b}
proof

```

    fix x
    assume  $x \in \{x \in A. f'(p_A x) = b\}$ 
    from this have  $x \in A f'(p_A x) = b$  by auto
    moreover from this  $\langle p_A \text{ permutes } A \rangle$  have  $p_A x \in A$  by (simp add:
permutes-in-image)
    moreover from  $\langle p_A \text{ permutes } A \rangle$  have  $x = \text{inv } p_A (p_A x)$ 
    using permutes-inverses(2) by fastforce
    ultimately show  $x \in \text{inv } p_A \{x \in A. f' x = b\}$  by auto
  qed
next
show  $\text{inv } p_A \{x \in A. f' x = b\} \subseteq \{x \in A. f'(p_A x) = b\}$ 
proof
  fix x
  assume  $x \in \text{inv } p_A \{x \in A. f' x = b\}$ 
  from this obtain  $x'$  where  $x: x = \text{inv } p_A x' x' \in A f' x' = b$  by auto
  from this  $\langle p_A \text{ permutes } A \rangle$  have  $x \in A$  by (simp add: permutes-in-image
permutes-inv)
  from  $\langle x = \text{inv } p_A x' \rangle \langle f' x' = b \rangle$  have  $f'(p_A x) = b$ 
  using  $\langle p_A \text{ permutes } A \rangle$  permutes-inverses(1) by fastforce
  from this  $\langle x \in A \rangle$  show  $x \in \{x \in A. f'(p_A x) = b\}$  by auto
  qed
qed
from this show ?thesis by blast
qed
also have  $\dots = (\lambda X. \text{inv } p_A X) \{(\lambda b. \{x \in A. f' x = b\}) B - \{\{\}\}\}$  by
auto
  finally show ?thesis by simp
qed
also have  $\dots = \text{image-mset } (\lambda X. \text{card } (\text{inv } p_A X)) (\text{mset-set } ((\lambda b. \{x \in A. f'
x = b\}) B - \{\{\}\}))$ 
  using  $\langle \text{inj-on } ((\cdot) (\text{inv } p_A)) ((\lambda b. \{x \in A. f' x = b\}) B - \{\{\}\}) \rangle$ 
  by (simp only: image-mset-mset-set[symmetric] image-mset.compositionality)
(meson comp-apply)
  also have  $\dots = \text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f' x = b\}) B - \{\{\}\}))$ 
  using  $\langle p_A \text{ permutes } A \rangle$  by (simp add: card-image inj-on-inv-into permutes-surj)
  finally show ?thesis .
qed

```

2.3.3 Existence of Permutation

lemma partition-implies-permutes:

assumes finite A

assumes partition-on A P partition-on A P'

assumes image-mset card (mset-set P') = image-mset card (mset-set P)

obtains p where p permutes A P' = $(\lambda X. p X) P$

proof –

from $\langle \text{partition-on } A P \rangle \langle \text{partition-on } A P' \rangle$ have finite P finite P'

using $\langle \text{finite } A \rangle$ finite-elements by blast+

from this image-mset card (mset-set P') = image-mset card (mset-set P)

```

obtain bij where bij-betw bij P P' and  $\forall X \in P. \text{card } X = \text{card } (\text{bij } X)$ 
  using image-mset-eq-implies-bij-betw by metis
have  $\forall X \in P. \exists p'. \text{bij-betw } p' X (\text{bij } X)$ 
proof
  fix X
  assume  $X \in P$ 
  from this have  $X \subseteq A$ 
    using  $\langle \text{partition-on } A \ P \rangle \text{ partition-onD1}$  by fastforce
  from this have finite X
    using  $\langle \text{finite } A \rangle \text{ rev-finite-subset}$  by blast
  from  $\langle X \in P \rangle$  have bij X  $\in P'$ 
    using  $\langle \text{bij-betw } \text{bij } P \ P' \rangle \text{ bij-betwE}$  by blast
  from this have bij X  $\subseteq A$ 
    using  $\langle \text{partition-on } A \ P' \rangle \text{ partition-onD1}$  by fastforce
  from this have finite (bij X)
    using  $\langle \text{finite } A \rangle \text{ rev-finite-subset}$  by blast
  from  $\langle X \in P \rangle$  have  $\text{card } X = \text{card } (\text{bij } X)$ 
    using  $\langle \forall X \in P. \text{card } X = \text{card } (\text{bij } X) \rangle$  by blast
  from this show  $\exists p'. \text{bij-betw } p' X (\text{bij } X)$ 
    using  $\langle \text{finite } (\text{bij } X) \rangle \langle \text{finite } X \rangle \text{ finite-same-card-bij}$  by blast
qed
from this have  $\exists p'. \forall X \in P. \text{bij-betw } (p' X) X (\text{bij } X)$  by metis
from this obtain p' where  $p': \forall X \in P. \text{bij-betw } (p' X) X (\text{bij } X) \dots$ 
define p where  $\bigwedge a. p \ a = (\text{if } a \in A \text{ then } p' (THE \ X. a \in X \wedge X \in P) \ a \text{ else}$ 
a)
have p permutes A
proof –
  have bij-betw p A A
proof –
  have disjoint-family-on bij P
proof
  fix X X'
  assume  $XX': X \in P \ X' \in P \ X \neq X'$ 
  from this have bij X  $\in P'$  bij X'  $\in P'$ 
    using  $\langle \text{bij-betw } \text{bij } P \ P' \rangle \text{ bij-betwE}$  by blast+
  moreover from  $XX'$  have bij X  $\neq \text{bij } X'$ 
    using  $\langle \text{bij-betw } \text{bij } P \ P' \rangle$  by (metis bij-betw-inv-into-left)
  ultimately show  $\text{bij } X \cap \text{bij } X' = \{\}$ 
    using  $\langle \text{partition-on } A \ P' \rangle$  by (meson partition-onE)
qed
moreover have bij-betw  $(\lambda a. p' (THE \ X. a \in X \wedge X \in P) \ a) X (\text{bij } X)$  if
 $X \in P$  for X
proof –
  from  $\langle X \in P \rangle$  have bij-betw  $(p' X) X (\text{bij } X)$ 
    using  $\langle \forall X \in P. \text{bij-betw } (p' X) X (\text{bij } X) \rangle$  by blast
  moreover from  $\langle X \in P \rangle$  have  $\forall a \in X. (THE \ X. a \in X \wedge X \in P) = X$ 
    using  $\langle \text{partition-on } A \ P \rangle \text{ partition-on-the-part-eq}$  by fastforce
  ultimately show ?thesis by (auto intro: bij-betw-congI)
qed

```

```

ultimately have bij-betw ( $\lambda a. p' (THE\ X. a \in X \wedge X \in P)\ a$ ) ( $\bigcup_{X \in P. X}$ )
( $\bigcup_{X \in P. bij\ X}$ )
  by (rule bij-betw-UNION-disjoint)
  moreover have ( $\bigcup_{X \in P. X} = A$ ) ( $\bigcup_{X \in P'. X} = A$ )
    using  $\langle partition-on\ A\ P \rangle \langle partition-on\ A\ P' \rangle$  partition-onD1 by auto
  moreover have ( $\bigcup_{X \in P. bij\ X} = (\bigcup_{X \in P'. X})$ )
    using  $\langle bij-betw\ bij\ P\ P' \rangle$  bij-betw-imp-surj-on by force
  ultimately have bij-betw ( $\lambda a. p' (THE\ X. a \in X \wedge X \in P)\ a$ )  $A\ A$  by simp
  moreover have  $\forall a \in A. p' (THE\ X. a \in X \wedge X \in P)\ a = p\ a$ 
    unfolding p-def by auto
  ultimately show ?thesis by (rule bij-betw-congI)
qed
moreover have  $p\ x = x$  if  $x \notin A$  for  $x$ 
  using  $\langle x \notin A \rangle$  p-def by auto
ultimately show ?thesis by (rule bij-imp-permutes)
qed
moreover have  $P' = (\lambda X. p\ ' X)\ ' P$ 
proof
  show  $P' \subseteq (\lambda X. p\ ' X)\ ' P$ 
  proof
    fix  $X$ 
    assume  $X \in P'$ 
    have in-P: the-inv-into  $P\ bij\ X \in P$ 
      using  $\langle X \in P' \rangle \langle bij-betw\ bij\ P\ P' \rangle$  bij-betwE bij-betw-the-inv-into by blast
    have eq-X: bij (the-inv-into  $P\ bij\ X$ ) =  $X$ 
      using  $\langle X \in P' \rangle \langle bij-betw\ bij\ P\ P' \rangle$ 
      by (meson f-the-inv-into-f-bij-betw)
    have  $X = p\ ' (the-inv-into\ P\ bij\ X)$ 
  proof
    from in-P have the-inv-into  $P\ bij\ X \subseteq A$ 
      using  $\langle partition-on\ A\ P \rangle$  partition-onD1 by fastforce
    have ( $\lambda a. p' (THE\ X. a \in X \wedge X \in P)\ a$ ) ' the-inv-into  $P\ bij\ X = X$ 
  proof
    show ( $\lambda a. p' (THE\ X. a \in X \wedge X \in P)\ a$ ) ' the-inv-into  $P\ bij\ X \subseteq X$ 
  proof
    fix  $x$ 
    assume  $x \in (\lambda a. p' (THE\ X. a \in X \wedge X \in P)\ a)$  ' the-inv-into  $P\ bij\ X$ 
    from this obtain  $a$  where a-in:  $a \in the-inv-into\ P\ bij\ X$ 
      and x-eq:  $x = p' (THE\ X. a \in X \wedge X \in P)\ a$  by blast
    have  $(THE\ X. a \in X \wedge X \in P) = the-inv-into\ P\ bij\ X$ 
      using a-in in-P  $\langle partition-on\ A\ P \rangle$  partition-on-the-part-eq
      by fastforce
    from this x-eq have x-eq:  $x = p' (the-inv-into\ P\ bij\ X)\ a$ 
      by auto
    from this have  $x \in bij\ (the-inv-into\ P\ bij\ X)$ 
      using a-in in-P bij-betwE  $p'$  by blast
    from this eq-X show  $x \in X$  by blast
  qed
qed
next

```

```

show  $X \subseteq (\lambda a. p' (THE X. a \in X \wedge X \in P) a)$  ‘ the-inv-into  $P$  bij  $X$ 
proof
  fix  $x$ 
  assume  $x \in X$ 
  let  $?X' = \text{the-inv-into } P \text{ bij } X$ 
  define  $x'$  where  $x' = \text{the-inv-into } ?X' (p' ?X') x$ 
  from in-P  $p' \text{ eq-} X$  have bij-betw: bij-betw  $(p' ?X') ?X' X$  by auto
  from bij-betw  $\langle x \in X \rangle$  have  $x' \in ?X'$ 
    unfolding  $x'\text{-def}$ 
    using bij-betwE bij-betw-the-inv-into by blast
  from this in-P have  $(THE X. x' \in X \wedge X \in P) = ?X'$ 
    using  $\langle \text{partition-on } A \ P \rangle$  partition-on-the-part-eq by fastforce
  from this  $\langle x \in X \rangle$  have  $x = p' (THE X. x' \in X \wedge X \in P) x'$ 
    unfolding  $x'\text{-def}$ 
    using bij-betw f-the-inv-into-f-bij-betw by fastforce
  from this  $\langle x' \in ?X' \rangle$  show  $x \in (\lambda a. p' (THE X. a \in X \wedge X \in P) a)$  ‘
the-inv-into  $P$  bij  $X$  ..
  qed
qed
from this  $\langle \text{the-inv-into } P \text{ bij } X \subseteq A \rangle$  show  $X \subseteq p$  ‘ the-inv-into  $P$  bij  $X$ 
  unfolding  $p\text{-def}$  by auto
next
show  $p$  ‘ the-inv-into  $P$  bij  $X \subseteq X$ 
proof
  fix  $x$ 
  assume  $x \in p$  ‘ the-inv-into  $P$  bij  $X$ 
  from this obtain  $x'$  where  $x = p x'$  and  $x' \in \text{the-inv-into } P \text{ bij } X$ 
    by auto
  have  $x' \in A$ 
    using  $\langle x' \in \text{the-inv-into } P \text{ bij } X \rangle$  assms(2) in-P partition-onD1 by
fastforce
  have eq:  $(THE X. x' \in X \wedge X \in P) = \text{the-inv-into } P \text{ bij } X$ 
    using  $\langle x' \in \text{the-inv-into } P \text{ bij } X \rangle$  assms(2) in-P partition-on-the-part-eq
by fastforce
  have  $p'$ :  $p' (\text{the-inv-into } P \text{ bij } X) x' \in X$ 
    using  $\langle x' \in \text{the-inv-into } P \text{ bij } X \rangle$  bij-betwE eq-X in-P  $p'$  by blast
  from  $\langle x = p x' \rangle \langle x' \in A \rangle$  eq  $p'$  show  $x \in X$ 
    unfolding  $p\text{-def}$  by auto
  qed
qed
moreover from  $\langle X \in P' \rangle \langle \text{bij-betw } P \ P' \rangle$  have the-inv-into  $P$  bij  $X \in P$ 
  using bij-betwE bij-betw-the-inv-into by blast
ultimately show  $X \in (\lambda X. p' X)$  ‘  $P$  ..
qed
next
show  $(\lambda X. p' X)$  ‘  $P \subseteq P'$ 
proof
  fix  $X'$ 
  assume  $X' \in (\lambda X. p' X)$  ‘  $P$ 

```

```

from this obtain  $X$  where  $X'\text{-eq}: X' = p \text{ ' } X$  and  $X \in P$  ..
from  $\langle X \in P \rangle$  have  $X \subseteq A$ 
  using assms(2) partition-onD1 by force
from  $\langle X \in P \rangle$   $p'$  have bij: bij-betw  $(p' X) X$  (bij  $X$ ) by auto
have  $p \text{ ' } X \in P'$ 
proof –
  from  $\langle X \in P \rangle$   $\langle \text{iij-betw } \text{iij } P P' \rangle$  have bij  $X \in P'$ 
    using bij-betwE by blast
  moreover have  $(\lambda a. p' (THE X. a \in X \wedge X \in P) a) \text{ ' } X = \text{iij } X$ 
  proof
    show  $(\lambda a. p' (THE X. a \in X \wedge X \in P) a) \text{ ' } X \subseteq \text{iij } X$ 
    proof
      fix  $x'$ 
      assume  $x' \in (\lambda a. p' (THE X. a \in X \wedge X \in P) a) \text{ ' } X$ 
      from this obtain  $x$  where  $x \in X$  and  $x'\text{-eq}: x' = p' (THE X. x \in X \wedge$ 
 $X \in P) x$  ..
      from  $\langle X \in P \rangle$   $\langle x \in X \rangle$  have eq-X:  $(THE X. x \in X \wedge X \in P) = X$ 
        using assms(2) partition-on-the-part-eq by fastforce
      from bij  $\langle x \in X \rangle$   $x'\text{-eq } eq\text{-}X$  show  $x' \in \text{iij } X$ 
        using bij-betwE by blast
      qed
    next
    show bij  $X \subseteq (\lambda a. p' (THE X. a \in X \wedge X \in P) a) \text{ ' } X$ 
    proof
      fix  $x'$ 
      assume  $x' \in \text{iij } X$ 
      let  $?x = \text{inv-into } X (p' X) x'$ 
      from  $\langle x' \in \text{iij } X \rangle$  bij have  $?x \in X$ 
        by (metis bij-betw-imp-surj-on inv-into-into)
      from this  $\langle X \in P \rangle$  have  $(THE X. ?x \in X \wedge X \in P) = X$ 
        using assms(2) partition-on-the-part-eq by fastforce
      from this  $\langle x' \in \text{iij } X \rangle$  bij have  $x' = p' (THE X. ?x \in X \wedge X \in P) ?x$ 
        using bij-betw-inv-into-right by fastforce
      moreover from  $\langle x' \in \text{iij } X \rangle$  bij have  $?x \in X$ 
        by (metis bij-betw-imp-surj-on inv-into-into)
      ultimately show  $x' \in (\lambda a. p' (THE X. a \in X \wedge X \in P) a) \text{ ' } X$  ..
      qed
    qed
    ultimately have  $(\lambda a. p' (THE X. a \in X \wedge X \in P) a) \text{ ' } X \in P'$  by simp
    have  $(\lambda a. p' (THE X. a \in X \wedge X \in P) a) \text{ ' } X = (\lambda a. \text{if } a \in A \text{ then } p'$ 
 $(THE X. a \in X \wedge X \in P) a \text{ else } a) \text{ ' } X$ 
    using  $\langle X \subseteq A \rangle$  by (auto intro: image-cong)
    from this show ?thesis
    using  $\langle (\lambda a. p' (THE X. a \in X \wedge X \in P) a) \text{ ' } X \in P' \rangle$  unfolding p-def
by auto
    qed
  from this  $X'\text{-eq}$  show  $X' \in P'$  by simp
  qed
qed

```


ultimately show *thesis* using that by *blast*
qed

lemma *permutes-domain-partition-eq*:

assumes $f \in A \rightarrow B$
assumes p_A permutes A
assumes $b \in B$
shows $p_A \text{ ' } \{x \in A. f\ x = b\} = \{x \in A. f\ (inv\ p_A\ x) = b\}$
proof
show $p_A \text{ ' } \{x \in A. f\ x = b\} \subseteq \{x \in A. f\ (inv\ p_A\ x) = b\}$
using $\langle p_A \text{ permutes } A \rangle$ *permutes-in-image permutes-inverses*(2) by *fastforce*
next
show $\{x \in A. f\ (inv\ p_A\ x) = b\} \subseteq p_A \text{ ' } \{x \in A. f\ x = b\}$
proof
fix x
assume $x \in \{x \in A. f\ (inv\ p_A\ x) = b\}$
from *this* have $x \in A$ $f\ (inv\ p_A\ x) = b$ by *auto*
from $\langle x \in A \rangle$ have $x = p_A\ (inv\ p_A\ x)$
using $\langle p_A \text{ permutes } A \rangle$ *permutes-inverses*(1) by *fastforce*
moreover from $\langle f\ (inv\ p_A\ x) = b \rangle \langle x \in A \rangle$ have $inv\ p_A\ x \in \{x \in A. f\ x = b\}$
by (*simp add: $\langle p_A \text{ permutes } A \rangle$ permutes-in-image permutes-inv*)
ultimately show $x \in p_A \text{ ' } \{x \in A. f\ x = b\}$..
qed
qed

lemma *image-domain-partition-eq*:

assumes $f \in A \rightarrow_E B$
assumes p_A permutes A
shows $(\lambda X. p_A \text{ ' } X) \text{ ' } ((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B) = (\lambda b. \{x \in A. f\ (inv\ p_A\ x) = b\}) \text{ ' } B$
proof
from $\langle f \in A \rightarrow_E B \rangle$ have $f \in A \rightarrow B$ by *auto*
note $eq = \text{permutes-domain-partition-eq}[OF\ \langle f \in A \rightarrow B \rangle\ \langle p_A \text{ permutes } A \rangle]$
show $(\lambda X. p_A \text{ ' } X) \text{ ' } ((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B) \subseteq (\lambda b. \{x \in A. f\ (inv\ p_A\ x) = b\}) \text{ ' } B$
proof
fix X
assume $X \in (\lambda X. p_A \text{ ' } X) \text{ ' } ((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B)$
from *this* obtain $b \in B$ and $X\text{-eq}$: $X = p_A \text{ ' } \{x \in A. f\ x = b\}$ by *auto*
from *this eq* have $X = \{x \in A. f\ (inv\ p_A\ x) = b\}$ by *simp*
from *this $\langle b \in B \rangle$* show $X \in (\lambda b. \{x \in A. f\ (inv\ p_A\ x) = b\}) \text{ ' } B$..
qed
next
from $\langle f \in A \rightarrow_E B \rangle$ have $f \in A \rightarrow B$ by *auto*
note $eq = \text{permutes-domain-partition-eq}[OF\ \langle f \in A \rightarrow B \rangle\ \langle p_A \text{ permutes } A \rangle,$
symmetric]
show $(\lambda b. \{x \in A. f\ (inv\ p_A\ x) = b\}) \text{ ' } B \subseteq (\lambda X. p_A \text{ ' } X) \text{ ' } ((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B)$
qed

```

proof
  fix  $X$ 
  assume  $X \in (\lambda b. \{x \in A. f (inv\ p_A\ x) = b\}) \text{ ' } B$ 
  from this obtain  $b \in B$  and  $X\text{-eq}: X = \{x \in A. f (inv\ p_A\ x) = b\}$ 
by auto
  from this eq have  $X = p_A \text{ ' } \{x \in A. f\ x = b\}$  by simp
  from this  $\langle b \in B \rangle$  show  $X \in (\lambda X. p_A \text{ ' } X) \text{ ' } (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B$  by
auto
qed
qed

```

lemma *multiset-of-partition-cards-eq-implies-permutes*:

```

assumes finite A finite B  $f \in A \rightarrow_E B$   $f' \in A \rightarrow_E B$ 
assumes eq: image-mset card (mset-set (( $\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}) =$ 
image-mset card (mset-set (( $\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B - \{\{\}\}))$ 
obtains  $p_A\ p_B$  where  $p_A$  permutes A  $p_B$  permutes B  $\forall x \in A. f\ x = p_B\ (f'\ (p_A\ x))$ 

```

proof –

```

have partition-on A (( $\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\})$ 
  using  $\langle f \in A \rightarrow_E B \rangle$  by (auto intro!: partition-onI)
moreover have partition-on A (( $\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B - \{\{\}\})$ 
  using  $\langle f' \in A \rightarrow_E B \rangle$  by (auto intro!: partition-onI)
moreover note partition-implyes-permutes[OF  $\langle \text{finite } A \rangle$  - - eq]
ultimately obtain  $p_A$  where  $p_A$  permutes A and
  inv-image-eq: (( $\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}) =$ 
  (( $\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B - \{\{\}\})$  by blast
from  $\langle p_A \text{ permutes } A \rangle$  have inj (( $\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B - \{\{\}\})$ 
  by (meson injI inj-image-eq-iff permutes-inj)
have inv-image-eq': (( $\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}) = (\lambda b. \{x \in A. f'\ (inv\ p_A\ x) = b\}) \text{ ' } B - \{\{\}\})$ 

```

proof –

```

note inv-image-eq
also have  $(\lambda X. p_A \text{ ' } X) \text{ ' } ((\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B - \{\{\}\}) = (\lambda b. \{x \in$ 
 $A. f'\ (inv\ p_A\ x) = b\}) \text{ ' } B - \{\{\}\}$ 
  using image-domain-partition-eq[OF  $\langle f' \in A \rightarrow_E B \rangle \langle p_A \text{ permutes } A \rangle$ 
  by (simp add: image-set-diff[OF  $\langle inj ((\lambda b. \{x \in A. f'\ x = b\}) \text{ ' } B - \{\{\}\}) \rangle$ 
finally show ?thesis .

```

qed

```

from  $\langle p_A \text{ permutes } A \rangle$  have inv p_A permutes A
  using permutes-inv by blast
have  $(\lambda x. f'\ (inv\ p_A\ x)) \in A \rightarrow_E B$ 
  using  $\langle f' \in A \rightarrow_E B \rangle \langle inv\ p_A \text{ permutes } A \rangle$  permutes-in-image by fastforce
from  $\langle f \in A \rightarrow_E B \rangle$  this  $\langle \text{finite } B \rangle$  obtain  $p_B$ 
  where  $p_B$  permutes B and  $eq'': \forall x \in A. f\ x = p_B\ (f'\ (inv\ p_A\ x))$ 
  using partitions-eq-implyes-permutes[OF  $\langle f \in A \rightarrow_E B \rangle \langle \text{finite } B \rangle$  - - inv-image-eq'] by blast
from  $\langle inv\ p_A \text{ permutes } A \rangle \langle p_B \text{ permutes } B \rangle eq''$  that show thesis by blast
qed

```

2.4 Bijections on Same Domain and Range

2.4.1 Existence of Domain Permutation

lemma *obtain-domain-permutation-for-two-bijections:*

assumes *bij-betw* f A B *bij-betw* f' A B

obtains p **where** p *permutes* A **and** $\forall a \in A. f\ a = f'\ (p\ a)$

proof –

let $?p = \lambda a. \text{if } a \in A \text{ then the-inv-into } A\ f'\ (f\ a) \text{ else } a$

have $?p$ *permutes* A

proof (*rule bij-imp-permutes*)

show *bij-betw* $?p$ A A

proof (*rule bij-betw-imageI*)

show *inj-on* $?p$ A

proof (*rule inj-onI*)

fix $a\ a'$

assume $a \in A\ a' \in A\ ?p\ a = ?p\ a'$

from *this* **have** *the-inv-into* $A\ f'\ (f\ a) = \text{the-inv-into } A\ f'\ (f\ a')$

using $\langle a \in A \rangle\ \langle a' \in A \rangle$ **by** *simp*

from *this* **have** $f\ a = f\ a'$

using $\langle a \in A \rangle\ \langle a' \in A \rangle$ *assms*

by (*metis bij-betwE f-the-inv-into-f-bij-betw*)

from *this* **show** $a = a'$

using $\langle a \in A \rangle\ \langle a' \in A \rangle$ *assms*

by (*metis bij-betw-inv-into-left*)

qed

next

show $?p\ 'A = A$

proof

show $?p\ 'A \subseteq A$

proof

fix a

assume $a \in ?p\ 'A$

from *this* **obtain** a' **where** $a' \in A$ **and** $a = \text{the-inv-into } A\ f'\ (f\ a')$ **by**

auto

from *this* *assms* **show** $a \in A$

by (*metis bij-betwE bij-betw-imp-inj-on bij-betw-imp-surj-on subset-iff the-inv-into-into*)

qed

next

show $A \subseteq ?p\ 'A$

proof

fix a

assume $a \in A$

from *this* *assms* **have** *the-inv-into* $A\ f\ (f'\ a) \in A$

by (*meson bij-betwE bij-betw-the-inv-into*)

moreover from $\langle a \in A \rangle$ *assms* **have** $a = \text{the-inv-into } A\ f'\ (f\ (\text{the-inv-into } A\ f\ (f'\ a)))$

by (*metis bij-betwE bij-betw-imp-inj-on f-the-inv-into-f-bij-betw the-inv-into-f-eq*)

ultimately show $a \in ?p\ 'A$ **by** *auto*

```

      qed
    qed
  qed
next
  fix a
  assume a  $\notin$  A
  from this show ?p a = a by auto
qed
moreover have  $\forall a \in A. f\ a = f'\ ( ?p\ a)$ 
  using  $\langle \text{bij-betw } f\ A\ B \rangle \langle \text{bij-betw } f'\ A\ B \rangle$ 
  using  $\text{bij-betwE } f\text{-the-inv-into-}f\text{-bij-betw}$  by fastforce
moreover note that
ultimately show thesis by auto
qed

```

2.4.2 Existence of Range Permutation

lemma *obtain-range-permutation-for-two-bijections:*

```

  assumes  $\text{bij-betw } f\ A\ B\ \text{bij-betw } f'\ A\ B$ 
  obtains p where p permutes B and  $\forall a \in A. f\ a = p\ (f'\ a)$ 
proof -
  let ?p =  $\lambda b. \text{if } b \in B \text{ then } f\ (\text{inv-into } A\ f'\ b) \text{ else } b$ 
  have ?p permutes B
  proof (rule  $\text{bij-imp-permutes}$ )
    show  $\text{bij-betw } ?p\ B\ B$ 
  proof (rule  $\text{bij-betw-imageI}$ )
    show  $\text{inj-on } ?p\ B$ 
  proof (rule  $\text{inj-onI}$ )
    fix b b'
    assume  $b \in B\ b' \in B\ ?p\ b = ?p\ b'$ 
    from this have  $f\ (\text{inv-into } A\ f'\ b) = f\ (\text{inv-into } A\ f'\ b')$ 
      using  $\langle b \in B \rangle \langle b' \in B \rangle$  by simp
    from this have  $\text{inv-into } A\ f'\ b = \text{inv-into } A\ f'\ b'$ 
      using  $\langle b \in B \rangle \langle b' \in B \rangle$  assms
    by (metis  $\text{bij-betw-imp-surj-on } \text{bij-betw-inv-into-left } \text{inv-into-into}$ )
    from this show  $b = b'$ 
      using  $\langle b \in B \rangle \langle b' \in B \rangle$  assms(2)
      by (metis  $\text{bij-betw-inv-into-right}$ )
  qed
  qed
next
  show ?p ' B = B
  proof
    from assms show  $?p\ ' B \subseteq B$ 
      by (auto simp add:  $\text{bij-betwE } \text{bij-betw-def } \text{inv-into-into}$ )
  next
    show  $B \subseteq ?p\ ' B$ 
  proof
    fix b
    assume  $b \in B$ 

```

```

    from this assms have  $f' (inv\text{-}into\ A\ f\ b) \in B$ 
    by (metis bij-betwE bij-betw-imp-surj-on inv-into-into)
    moreover have  $b = ?p\ (f' (inv\text{-}into\ A\ f\ b))$ 
    using assms  $\langle f' (inv\text{-}into\ A\ f\ b) \in B \rangle \langle b \in B \rangle$ 
    by (auto simp add: bij-betw-imp-surj-on bij-betw-inv-into-left bij-betw-inv-into-right
inv-into-into)
    ultimately show  $b \in ?p\ `B$  by auto
  qed
qed
qed
next
fix b
assume  $b \notin B$ 
from this show  $?p\ b = b$  by auto
qed
moreover have  $\forall a \in A. f\ a = ?p\ (f'\ a)$ 
using  $\langle bij\text{-}betw\ f'\ A\ B \rangle$  bij-betw-inv-into-left bij-betwE by fastforce
moreover note that
ultimately show thesis by auto
qed
end

```

3 Definition of Equivalence Classes

```

theory Equiv-Relations-on-Functions
imports
  Preliminaries
  Twelffold-Way-Core
begin

```

3.1 Permutation on the Domain

definition *domain-permutation*

where

$domain\text{-}permutation\ A\ B = \{(f, f') \in (A \rightarrow_E B) \times (A \rightarrow_E B). \exists p. p\ \text{permutes}\ A \wedge (\forall x \in A. f\ x = f'\ (p\ x))\}$

lemma *equiv-domain-permutation:*

$equiv\ (A \rightarrow_E B)\ (domain\text{-}permutation\ A\ B)$

proof (rule equivI)

show $domain\text{-}permutation\ A\ B \subseteq (A \rightarrow_E B) \times (A \rightarrow_E B)$

unfolding *domain-permutation-def* **by** auto

next

show $refl\text{-}on\ (A \rightarrow_E B)\ (domain\text{-}permutation\ A\ B)$

proof (rule refl-onI)

fix f

assume $f \in A \rightarrow_E B$

from this show $(f, f) \in domain\text{-}permutation\ A\ B$

```

    using permutes-id unfolding domain-permutation-def by fastforce
  qed
next
show sym (domain-permutation A B)
proof (rule symI)
  fix f f'
  assume (f, f') ∈ domain-permutation A B
  from this obtain p where p permutes A and  $\forall x \in A. f x = f' (p x)$ 
    unfolding domain-permutation-def by auto
  from  $\langle (f, f') \in \text{domain-permutation } A \ B \rangle$  have  $f \in A \rightarrow_E B$   $f' \in A \rightarrow_E B$ 
    unfolding domain-permutation-def by auto
  moreover from  $\langle p \text{ permutes } A \rangle$  have  $\text{inv } p \text{ permutes } A$ 
    by (simp add: permutes-inv)
  moreover from  $\langle p \text{ permutes } A \rangle \langle \forall x \in A. f x = f' (p x) \rangle$  have  $\forall x \in A. f' x = f$ 
    (inv p x)
    using permutes-in-image permutes-inverses(1) by (metis (mono-tags, opaque-lifting))
  ultimately show (f', f) ∈ domain-permutation A B
    unfolding domain-permutation-def by auto
  qed
next
show trans (domain-permutation A B)
proof (rule transI)
  fix f f' f''
  assume (f, f') ∈ domain-permutation A B (f', f'') ∈ domain-permutation A B
  from  $\langle (f, f') \in \rightarrow \rangle$  obtain p where p permutes A and  $\forall x \in A. f x = f' (p x)$ 
    unfolding domain-permutation-def by auto
  from  $\langle (f', f'') \in \rightarrow \rangle$  obtain p' where p' permutes A and  $\forall x \in A. f' x = f'' (p' x)$ 
    (inv p' x)
    unfolding domain-permutation-def by auto
  from  $\langle (f, f') \in \text{domain-permutation } A \ B \rangle$  have  $f \in A \rightarrow_E B$ 
    unfolding domain-permutation-def by auto
  moreover from  $\langle (f', f'') \in \text{domain-permutation } A \ B \rangle$  have  $f'' \in A \rightarrow_E B$ 
    unfolding domain-permutation-def by auto
  moreover from  $\langle p \text{ permutes } A \rangle \langle p' \text{ permutes } A \rangle$  have  $(p' \circ p) \text{ permutes } A$ 
    by (simp add: permutes-compose)
  moreover have  $\forall x \in A. f x = f'' ((p' \circ p) x)$ 
    using  $\langle \forall x \in A. f x = f' (p x) \rangle \langle \forall x \in A. f' x = f'' (p' x) \rangle \langle p \text{ permutes } A \rangle$ 
    by (simp add: permutes-in-image)
  ultimately show (f, f'') ∈ domain-permutation A B
    unfolding domain-permutation-def by auto
  qed
qed

```

3.1.1 Respecting Functions

lemma *inj-on-respects-domain-permutation:*

$(\lambda f. \text{inj-on } f \ A) \text{ respects domain-permutation } A \ B$

proof (rule congruentI)

fix f f'

```

assume (f, f') ∈ domain-permutation A B
from this obtain p where p: p permutes A ∀ x ∈ A. f x = f' (p x)
  unfolding domain-permutation-def by auto
have inv-p: ∀ x ∈ A. f' x = f (inv p x)
  using p by (metis permutes-inverses(1) permutes-not-in)
show inj-on f A ⟷ inj-on f' A
proof
  assume inj-on f A
  show inj-on f' A
  proof (rule inj-onI)
    fix a a'
    assume a ∈ A a' ∈ A f' a = f' a'
    from this ⟨p permutes A⟩ have inv p a ∈ A inv p a' ∈ A
      by (simp add: permutes-in-image permutes-inv)+
    have f (inv p a) = f (inv p a')
      using ⟨f' a = f' a'⟩ ⟨a ∈ A⟩ ⟨a' ∈ A⟩ inv-p by auto
    from ⟨inj-on f A⟩ this ⟨inv p a ∈ A⟩ ⟨inv p a' ∈ A⟩ have inv p a = inv p a'
      using inj-on-contrad by fastforce
    from this show a = a'
      by (metis ⟨p permutes A⟩ permutes-inverses(1))
  qed
next
  assume inj-on f' A
  from this p show inj-on f A
    unfolding inj-on-def
    by (metis inj-on-contrad permutes-in-image permutes-inj-on)
  qed
qed

lemma image-respects-domain-permutation:
  (λf. f ' A) respects (domain-permutation A B)
proof (rule congruentI)
  fix f f'
  assume (f, f') ∈ domain-permutation A B
  from this obtain p where p: p permutes A and f-eq: ∀ x ∈ A. f x = f' (p x)
    unfolding domain-permutation-def by auto
  show f ' A = f' ' A
  proof
    from p f-eq show f ' A ⊆ f' ' A
      by (auto simp add: permutes-in-image)
  next
    from ⟨p permutes A⟩ ⟨∀ x ∈ A. f x = f' (p x)⟩ have ∀ x ∈ A. f' x = f (inv p x)
      using permutes-in-image permutes-inverses(1) by (metis (mono-tags, opaque-lifting))
    from this show f' ' A ⊆ f ' A
      using ⟨p permutes A⟩ by (auto simp add: permutes-inv permutes-in-image)
  qed
qed

```

lemma surjective-respects-domain-permutation:

($\lambda f. f \text{ ` } A = B$) *respects domain-permutation* $A B$
by (*metis image-respects-domain-permutation congruentD congruentI*)

lemma *bij-betw-respects-domain-permutation:*

($\lambda f. \text{bij-betw } f A B$) *respects domain-permutation* $A B$

proof (*rule congruentI*)

fix $f f'$

assume $(f, f') \in \text{domain-permutation } A B$

from this obtain p **where** p *permutes* A **and** $\forall x \in A. f x = f' (p x)$

unfolding *domain-permutation-def* **by** *auto*

have $\text{bij-betw } f A B \longleftrightarrow \text{bij-betw } (f' \circ p) A B$

using $\langle \forall x \in A. f x = f' (p x) \rangle$

by (*metis (mono-tags, opaque-lifting) comp-apply[of f' p] bij-betw-cong[of A f f' \circ p B]*)

also have $\dots \longleftrightarrow \text{bij-betw } f' A B$

using $\langle p \text{ permutes } A \rangle$

by (*auto intro!: bij-betw-comp-iff[symmetric] permutes-imp-bij*)

finally show $\text{bij-betw } f A B \longleftrightarrow \text{bij-betw } f' A B$.

qed

lemma *image-mset-respects-domain-permutation:*

shows ($\lambda f. \text{image-mset } f (\text{mset-set } A)$) *respects (domain-permutation* $A B$)

proof (*rule congruentI*)

fix $f f'$

assume $(f, f') \in \text{domain-permutation } A B$

from this obtain p **where** p *permutes* A **and** $\forall x \in A. f x = f' (p x)$

unfolding *domain-permutation-def* **by** *auto*

from this show $\text{image-mset } f (\text{mset-set } A) = \text{image-mset } f' (\text{mset-set } A)$

using *permutes-implies-image-mset-eq* **by** *fastforce*

qed

3.2 Permutation on the Range

definition *range-permutation*

where

$\text{range-permutation } A B = \{(f, f') \in (A \rightarrow_E B) \times (A \rightarrow_E B). \exists p. p \text{ permutes } B$
 $\wedge (\forall x \in A. f x = p (f' x))\}$

lemma *equiv-range-permutation:*

equiv $(A \rightarrow_E B)$ (*range-permutation* $A B$)

proof (*rule equivI*)

show $\text{range-permutation } A B \subseteq (A \rightarrow_E B) \times (A \rightarrow_E B)$

unfolding *range-permutation-def* **by** *auto*

next

show *refl-on* $(A \rightarrow_E B)$ (*range-permutation* $A B$)

proof (*rule refl-onI*)

fix f

assume $f \in A \rightarrow_E B$

from this show $(f, f) \in \text{range-permutation } A B$


```

    using permutes-id unfolding range-permutation-def by fastforce
  qed
next
show sym (range-permutation A B)
proof (rule symI)
  fix f f'
  assume (f, f') ∈ range-permutation A B
  from this obtain p where p permutes B and  $\forall x \in A. f x = p (f' x)$ 
    unfolding range-permutation-def by auto
  from  $\langle (f, f') \in \text{range-permutation } A \ B \rangle$  have  $f \in A \rightarrow_E B \ f' \in A \rightarrow_E B$ 
    unfolding range-permutation-def by auto
  moreover from  $\langle p \text{ permutes } B \rangle$  have inv p permutes B
    by (simp add: permutes-inv)
  moreover from  $\langle p \text{ permutes } B \rangle \ \langle \forall x \in A. f x = p (f' x) \rangle$  have  $\forall x \in A. f' x =$ 
    inv p (f x)
    by (simp add: permutes-inverses(2))
  ultimately show (f', f) ∈ range-permutation A B
    unfolding range-permutation-def by auto
  qed
next
show trans (range-permutation A B)
proof (rule transI)
  fix f f' f''
  assume (f, f') ∈ range-permutation A B (f', f'') ∈ range-permutation A B
  from  $\langle (f, f') \in \rightarrow \rangle$  obtain p where p permutes B and  $\forall x \in A. f x = p (f' x)$ 
    unfolding range-permutation-def by auto
  from  $\langle (f', f'') \in \rightarrow \rangle$  obtain p' where p' permutes B and  $\forall x \in A. f' x = p' (f'' x)$ 
    unfolding range-permutation-def by auto
  moreover from  $\langle p \text{ permutes } B \rangle \ \langle p' \text{ permutes } B \rangle$  have  $(p \circ p') \text{ permutes } B$ 
    by (simp add: permutes-compose)
  moreover have  $\forall x \in A. f x = (p \circ p') (f'' x)$ 
    using  $\langle \forall x \in A. f x = p (f' x) \rangle \ \langle \forall x \in A. f' x = p' (f'' x) \rangle$  by auto
  ultimately show (f, f'') ∈ range-permutation A B
    unfolding range-permutation-def by auto
  qed
qed

```

3.2.1 Respecting Functions

lemma *inj-on-respects-range-permutation:*
 $(\lambda f. \text{inj-on } f \ A) \text{ respects range-permutation } A \ B$
proof (rule congruentI)
 fix f f'
 assume (f, f') ∈ range-permutation A B

```

from this obtain  $p$  where  $p$ :  $p$  permutes  $B \forall x \in A. f x = p (f' x)$ 
  unfolding range-permutation-def by auto
have  $inv\text{-}p$ :  $\forall x \in A. f' x = inv\ p (f x)$ 
  using  $p$  by (simp add: permutes-inverses(2))
show  $inj\text{-}on\ f\ A \longleftrightarrow inj\text{-}on\ f'\ A$ 
proof
  assume  $inj\text{-}on\ f\ A$ 
  from this  $p$  show  $inj\text{-}on\ f'\ A$ 
    unfolding inj-on-def by auto
next
  assume  $inj\text{-}on\ f'\ A$ 
  from this  $inv\text{-}p$  show  $inj\text{-}on\ f\ A$ 
    unfolding inj-on-def by auto
qed
qed

lemma surj-on-respects-range-permutation:
   $(\lambda f. f \text{ ' } A = B)$  respects range-permutation  $A\ B$ 
proof (rule congruentI)
  fix  $f\ f'$ 
  assume  $a$ :  $(f, f') \in range\text{-}permutation\ A\ B$ 
  from this have  $f \in A \rightarrow_E B\ f' \in A \rightarrow_E B$ 
    unfolding range-permutation-def by auto
  from  $a$  obtain  $p$  where  $p$ :  $p$  permutes  $B \forall x \in A. f x = p (f' x)$ 
    unfolding range-permutation-def by auto
  have  $1$ :  $f \text{ ' } A = (\lambda x. p (f' x)) \text{ ' } A$ 
    using  $p$  by (meson image-cong)
  have  $2$ :  $inv\ p \text{ ' } ((\lambda x. p (f' x)) \text{ ' } A) = f' \text{ ' } A$ 
    using  $p$  by (simp add: image-image image-inv-f-f permutes-inj)
  show  $(f \text{ ' } A = B) = (f' \text{ ' } A = B)$ 
  proof
    assume  $f \text{ ' } A = B$ 
    from this  $1\ 2$  show  $f' \text{ ' } A = B$ 
      using  $p$  by (simp add: permutes-image permutes-inv)
  next
    assume  $f' \text{ ' } A = B$ 
    from this  $1\ 2$  show  $f \text{ ' } A = B$ 
      using  $p$  by (metis image-image permutes-image)
  qed
qed

lemma bij-betw-respects-range-permutation:
   $(\lambda f. bij\text{-}betw\ f\ A\ B)$  respects range-permutation  $A\ B$ 
proof (rule congruentI)
  fix  $f\ f'$ 
  assume  $(f, f') \in range\text{-}permutation\ A\ B$ 
  from this obtain  $p$  where  $p$  permutes  $B$  and  $\forall x \in A. f x = p (f' x)$ 
    and  $f' \in A \rightarrow_E B$ 
    unfolding range-permutation-def by auto

```

have $\text{bij-betw } f \ A \ B \longleftrightarrow \text{bij-betw } (p \circ f') \ A \ B$
using $\langle \forall x \in A. f \ x = p \ (f' \ x) \rangle$
by (*metis* (*mono-tags*, *opaque-lifting*) *bij-betw-cong comp-apply*)
also have $\dots \longleftrightarrow \text{bij-betw } f' \ A \ B$
using $\langle f' \in A \rightarrow_E B \rangle \langle p \text{ permutes } B \rangle$
by (*auto intro!*: *bij-betw-comp-iff2[symmetric]* *permutes-imp-bij*)
finally show $\text{bij-betw } f \ A \ B \longleftrightarrow \text{bij-betw } f' \ A \ B$.
qed

lemma *domain-partitions-respects-range-permutation*:
 $(\lambda f. (\lambda b. \{x \in A. f \ x = b\}) \text{ ` } B - \{\{\}\}) \text{ respects range-permutation } A \ B$
proof (*rule congruentI*)
fix $f \ f'$
assume $(f, f') \in \text{range-permutation } A \ B$
from this obtain p **where** p : $p \text{ permutes } B \ \forall x \in A. f \ x = p \ (f' \ x)$
unfolding *range-permutation-def* **by** *blast*
have $\{\} \in (\lambda b. \{x \in A. f' \ x = b\}) \text{ ` } B \longleftrightarrow \neg (\forall b \in B. \exists x \in A. f' \ x = b)$ **by** *auto*
also have $(\forall b \in B. \exists x \in A. f' \ x = b) \longleftrightarrow (\forall b \in B. \exists x \in A. p \ (f' \ x) = b)$
proof
assume $\forall b \in B. \exists x \in A. f' \ x = b$
from this show $\forall b \in B. \exists x \in A. p \ (f' \ x) = b$
using $\langle p \text{ permutes } B \rangle$ **unfolding** *permutes-def* **by** *metis*
next
assume $\forall b \in B. \exists x \in A. p \ (f' \ x) = b$
from this show $\forall b \in B. \exists x \in A. f' \ x = b$
using $\langle p \text{ permutes } B \rangle$ **by** (*metis* *bij-betwE permutes-imp-bij permutes-inverses(2)*)
qed
also have $\neg (\forall b \in B. \exists x \in A. p \ (f' \ x) = b) \longleftrightarrow \{\} \in (\lambda b. \{x \in A. p \ (f' \ x) = b\})$
 $\text{ ` } B$ **by** *auto*
finally have $\{\} \in (\lambda b. \{x \in A. f' \ x = b\}) \text{ ` } B \longleftrightarrow \{\} \in (\lambda b. \{x \in A. p \ (f' \ x) = b\}) \text{ ` } B$.
moreover have $(\lambda b. \{x \in A. f' \ x = b\}) \text{ ` } B = (\lambda b. \{x \in A. p \ (f' \ x) = b\}) \text{ ` } B$
using $\langle p \text{ permutes } B \rangle$ *permutes-implies-inv-image-on-eq* **by** *blast*
ultimately have $(\lambda b. \{x \in A. f' \ x = b\}) \text{ ` } B - \{\{\}\} = (\lambda b. \{x \in A. p \ (f' \ x) = b\}) \text{ ` } B - \{\{\}\}$ **by** *auto*
also have $\dots = (\lambda b. \{x \in A. f \ x = b\}) \text{ ` } B - \{\{\}\}$
using $\langle \forall x \in A. f \ x = p \ (f' \ x) \rangle$ *Collect-cong image-cong* **by** *auto*
finally show $(\lambda b. \{x \in A. f \ x = b\}) \text{ ` } B - \{\{\}\} = (\lambda b. \{x \in A. f' \ x = b\}) \text{ ` } B - \{\{\}\}$..
qed

3.3 Permutation on the Domain and the Range

definition *domain-and-range-permutation*

where

$\text{domain-and-range-permutation } A \ B = \{(f, f') \in (A \rightarrow_E B) \times (A \rightarrow_E B).$
 $\exists p_A \ p_B. p_A \text{ permutes } A \wedge p_B \text{ permutes } B \wedge (\forall x \in A. f \ x = p_B \ (f' \ (p_A \ x)))\}$

```

lemma equiv-domain-and-range-permutation:
  equiv (A →E B) (domain-and-range-permutation A B)
proof (rule equivI)
  show domain-and-range-permutation A B ⊆ (A →E B) × (A →E B)
    unfolding domain-and-range-permutation-def by auto
next
  show refl-on (A →E B) (domain-and-range-permutation A B)
proof (rule refl-onI)
  fix f
  assume f ∈ A →E B
  from this show (f, f) ∈ domain-and-range-permutation A B
    using permutes-id[of A] permutes-id[of B]
    unfolding domain-and-range-permutation-def by fastforce
  qed
next
  show sym (domain-and-range-permutation A B)
proof (rule symI)
  fix f f'
  assume (f, f') ∈ domain-and-range-permutation A B
  from this obtain pA pB where pA permutes A pB permutes B and ∀ x ∈ A. f
x = pB (f' (pA x))
    unfolding domain-and-range-permutation-def by auto
  from ⟨(f, f') ∈ domain-and-range-permutation A B⟩ have f: f ∈ A →E B f'
∈ A →E B
    unfolding domain-and-range-permutation-def by auto
  moreover from ⟨pA permutes A⟩ ⟨pB permutes B⟩ have inv pA permutes A
inv pB permutes B
    by (auto simp add: permutes-inv)
  moreover from ⟨∀ x ∈ A. f x = pB (f' (pA x))⟩ have ∀ x ∈ A. f' x = inv pB (f
(inv pA x))
    using ⟨pA permutes A⟩ ⟨pB permutes B⟩ ⟨inv pA permutes A⟩ ⟨inv pB permutes
B⟩
    by (metis (no-types, lifting) bij-betwE bij-inv-eq-iff permutes-bij permutes-imp-bij)
  ultimately show (f', f) ∈ domain-and-range-permutation A B
    unfolding domain-and-range-permutation-def by auto
  qed
next
  show trans (domain-and-range-permutation A B)
proof (rule transI)
  fix f f' f''
  assume (f, f') ∈ domain-and-range-permutation A B
  assume (f', f'') ∈ domain-and-range-permutation A B
  from ⟨(f, f') ∈ →⟩ obtain pA pB where
    pA permutes A pB permutes B and ∀ x ∈ A. f x = pB (f' (pA x))
    unfolding domain-and-range-permutation-def by auto
  from ⟨(f', f'') ∈ →⟩ obtain p'A p'B where
    p'A permutes A p'B permutes B and ∀ x ∈ A. f' x = p'B (f'' (p'A x))
    unfolding domain-and-range-permutation-def by auto
  from ⟨(f, f') ∈ domain-and-range-permutation A B⟩ have f ∈ A →E B

```

unfolding *domain-and-range-permutation-def* **by** *auto*
moreover from $\langle f', f'' \rangle \in \text{domain-and-range-permutation } A \ B \rangle$ **have** $f'' \in A \rightarrow_E B$
unfolding *domain-and-range-permutation-def* **by** *auto*
moreover from $\langle p_A \text{ permutes } A \rangle \langle p'_A \text{ permutes } A \rangle$ **have** $(p'_A \circ p_A) \text{ permutes } A$
by (*simp add: permutes-compose*)
moreover from $\langle p_B \text{ permutes } B \rangle \langle p'_B \text{ permutes } B \rangle$ **have** $(p_B \circ p'_B) \text{ permutes } B$
by (*simp add: permutes-compose*)
moreover have $\forall x \in A. f \ x = (p_B \circ p'_B) (f'' ((p'_A \circ p_A) \ x))$
using $\langle \forall x \in A. f' \ x = p'_B (f'' (p'_A \ x)) \rangle \langle \forall x \in A. f \ x = p_B (f' (p_A \ x)) \rangle \langle p_A \text{ permutes } A \rangle$
by (*simp add: permutes-in-image*)
ultimately show $(f, f'') \in \text{domain-and-range-permutation } A \ B$
unfolding *domain-and-range-permutation-def* **by** *fastforce*
qed
qed

3.3.1 Respecting Functions

lemma *inj-on-respects-domain-and-range-permutation:*

$(\lambda f. \text{inj-on } f \ A) \text{ respects domain-and-range-permutation } A \ B$

proof (*rule congruentI*)

fix $f \ f'$

assume $(f, f') \in \text{domain-and-range-permutation } A \ B$

from this obtain $p_A \ p_B$ **where** $p_A \text{ permutes } A \ p_B \text{ permutes } B$ **and** $\forall x \in A. f \ x = p_B (f' (p_A \ x))$

unfolding *domain-and-range-permutation-def* **by** *auto*

from $\langle (f, f') \in \text{domain-and-range-permutation } A \ B \rangle$ **have** $f' \restriction A \subseteq B$

unfolding *domain-and-range-permutation-def* **by** *auto*

from $\langle p_A \text{ permutes } A \rangle$ **have** $p_A \restriction A = A$ **by** (*auto simp add: permutes-image*)

from $\langle p_A \text{ permutes } A \rangle$ **have** $\text{inj-on } p_A \ A$

using *bij-betw-imp-inj-on permutes-imp-bij* **by** *blast*

from $\langle p_B \text{ permutes } B \rangle$ **have** $\text{inj-on } p_B \ B$

using *bij-betw-imp-inj-on permutes-imp-bij* **by** *blast*

show $\text{inj-on } f \ A \longleftrightarrow \text{inj-on } f' \ A$

proof —

have $\text{inj-on } f \ A \longleftrightarrow \text{inj-on } (\lambda x. p_B (f' (p_A \ x))) \ A$

using $\langle \forall x \in A. f \ x = p_B (f' (p_A \ x)) \rangle$ *inj-on-cong comp-apply* **by** *fastforce*

have $\text{inj-on } f \ A \longleftrightarrow \text{inj-on } (p_B \circ f' \circ p_A) \ A$

by (*simp add: $\langle \forall x \in A. f \ x = p_B (f' (p_A \ x)) \rangle$ inj-on-def*)

also have $\text{inj-on } (p_B \circ f' \circ p_A) \ A \longleftrightarrow \text{inj-on } (p_B \circ f') \ A$

using $\langle \text{inj-on } p_A \ A \rangle \langle p_A \restriction A = A \rangle$

by (*auto dest: inj-on-imageI intro: comp-inj-on*)

also have $\text{inj-on } (p_B \circ f') \ A \longleftrightarrow \text{inj-on } f' \ A$

using $\langle \text{inj-on } p_B \ B \rangle \langle f' \restriction A \subseteq B \rangle$

by (*auto dest: inj-on-imageI2 intro: comp-inj-on inj-on-subset*)

finally show *?thesis* .

qed
qed

lemma *surjective-respects-domain-and-range-permutation:*

$(\lambda f. f \text{ ' } A = B)$ respects domain-and-range-permutation $A \ B$

proof (rule congruentI)

fix $f f'$

assume $(f, f') \in \text{domain-and-range-permutation } A \ B$

from this obtain $p_A \ p_B$ where

$\text{permutes } p_A \text{ permutes } A \ p_B \text{ permutes } B$ and $\forall x \in A. f \ x = p_B \ (f' \ (p_A \ x))$

unfolding domain-and-range-permutation-def by auto

from permutes have $p_A \text{ ' } A = A \ p_B \text{ ' } B = B$ by (auto simp add: permutes-image)

from $\langle p_B \text{ permutes } B \rangle$ have $\text{inj } p_B$ by (simp add: permutes-inj)

show $(f \text{ ' } A = B) \longleftrightarrow (f' \text{ ' } A = B)$

proof –

have $f \text{ ' } A = B \longleftrightarrow (\lambda x. p_B \ (f' \ (p_A \ x))) \text{ ' } A = B$

using $\langle \forall x \in A. f \ x = p_B \ (f' \ (p_A \ x)) \rangle$ by (metis (mono-tags, lifting) image-cong)

also have $(\lambda x. p_B \ (f' \ (p_A \ x))) \text{ ' } A = B \longleftrightarrow (\lambda x. p_B \ (f' \ x)) \text{ ' } A = B$

using $\langle p_A \text{ ' } A = A \rangle$ by (metis image-image)

also have $(\lambda x. p_B \ (f' \ x)) \text{ ' } A = B \longleftrightarrow (f' \text{ ' } A = B)$

using $\langle p_B \text{ ' } B = B \rangle \langle \text{inj } p_B \rangle$ by (metis image-image image-inv-f-f)

finally show ?thesis .

qed

qed

lemma *bij-betw-respects-domain-and-range-permutation:*

$(\lambda f. \text{bij-betw } f \ A \ B)$ respects domain-and-range-permutation $A \ B$

proof (rule congruentI)

fix $f f'$

assume $(f, f') \in \text{domain-and-range-permutation } A \ B$

from this obtain $p_A \ p_B$ where $p_A \text{ permutes } A \ p_B \text{ permutes } B$

and $\forall x \in A. f \ x = p_B \ (f' \ (p_A \ x))$ and $f' \in A \rightarrow_E \ B$

unfolding domain-and-range-permutation-def by auto

have $\text{bij-betw } f \ A \ B \longleftrightarrow \text{bij-betw } (p_B \circ f' \circ p_A) \ A \ B$

using $\langle \forall x \in A. f \ x = p_B \ (f' \ (p_A \ x)) \rangle \text{bij-betw-congI}$ by fastforce

also have $\dots \longleftrightarrow \text{bij-betw } (p_B \circ f') \ A \ B$

using $\langle p_A \text{ permutes } A \rangle$

by (auto intro!: bij-betw-comp-iff[symmetric] permutes-imp-bij)

also have $\dots \longleftrightarrow \text{bij-betw } f' \ A \ B$

using $\langle f' \in A \rightarrow_E \ B \rangle \langle p_B \text{ permutes } B \rangle$

by (auto intro!: bij-betw-comp-iff2[symmetric] permutes-imp-bij)

finally show $\text{bij-betw } f \ A \ B \longleftrightarrow \text{bij-betw } f' \ A \ B$.

qed

lemma *count-image-mset':*

$\text{count } (\text{image-mset } f \ A) \ x = \text{sum } (\text{count } A) \ \{x' \in \text{set-mset } A. f \ x' = x\}$

proof –

have $\text{count } (\text{image-mset } f \ A) \ x = \text{sum } (\text{count } A) \ (f \text{ - ' } \{x\} \cap \text{set-mset } A)$

unfolding count-image-mset ..

also have $\dots = \text{sum } (\text{count } A) \{x' \in \text{set-mset } A. f x' = x\}$
proof –
have $(f - \{x\} \cap \text{set-mset } A) = \{x' \in \text{set-mset } A. f x' = x\}$ **by** *blast*
from this show *?thesis by simp*
qed
finally show *?thesis* .
qed

lemma *multiset-of-partition-cards-respects-domain-and-range-permutation:*

assumes *finite B*
shows $(\lambda f. \text{image-mset } (\lambda X. \text{card } X) (\text{mset-set } ((\lambda b. \{x \in A. f x = b\}) ' B - \{\{\}\})))$ *respects domain-and-range-permutation A B*
proof (*rule congruentI*)
fix $f f'$
assume $(f, f') \in \text{domain-and-range-permutation } A B$
from this obtain $p_A p_B$ **where** p_A *permutes A* p_B *permutes B* $\forall x \in A. f x = p_B (f' (p_A x))$
unfolding *domain-and-range-permutation-def* **by** *auto*
have $(\lambda b. \{x \in A. f x = b\}) ' B = (\lambda b. \{x \in A. p_B (f' (p_A x)) = b\}) ' B$
using $\langle \forall x \in A. f x = p_B (f' (p_A x)) \rangle$ **by** *auto*
from this have $\text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f x = b\}) ' B - \{\{\}\}))$
 $=$
 $\text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. p_B (f' (p_A x)) = b\}) ' B - \{\{\}\}))$ **by** *simp*
also have $\text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. p_B (f' (p_A x)) = b\}) ' B - \{\{\}\})) =$
 $\text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f' (p_A x) = b\}) ' B - \{\{\}\}))$
using *permutes-implies-inv-image-on-eq[OF $\langle p_B \text{ permutes } B \rangle$, of A]* **by** *metis*
also have $\text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f' (p_A x) = b\}) ' B - \{\{\}\}))$
 $=$
 $\text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f' x = b\}) ' B - \{\{\}\}))$
proof (*rule multiset-eqI*)
fix n
have $\text{bij-betw } (\lambda X. p_A ' X) \{X \in (\lambda b. \{x \in A. f' (p_A x) = b\}) ' B - \{\{\}\}. \text{card } X = n\} \{X \in (\lambda b. \{x \in A. f' x = b\}) ' B - \{\{\}\}. \text{card } X = n\}$
proof (*rule bij-betw-byWitness*)
show $\forall X \in \{X \in (\lambda b. \{x \in A. f' (p_A x) = b\}) ' B - \{\{\}\}. \text{card } X = n\}. \text{inv } p_A ' p_A ' X = X$
by (*meson $\langle p_A \text{ permutes } A \rangle \text{ image-inv-f-f permutes-inj}$*)
show $\forall X \in \{X \in (\lambda b. \{x \in A. f' x = b\}) ' B - \{\{\}\}. \text{card } X = n\}. p_A ' \text{inv } p_A ' X = X$
by (*meson $\langle p_A \text{ permutes } A \rangle \text{ image-f-inv-f permutes-surj}$*)
show $(\lambda X. p_A ' X) ' \{X \in (\lambda b. \{x \in A. f' (p_A x) = b\}) ' B - \{\{\}\}. \text{card } X = n\} \subseteq \{X \in (\lambda b. \{x \in A. f' x = b\}) ' B - \{\{\}\}. \text{card } X = n\}$
proof –
have $\text{card } (p_A ' \{x \in A. f' (p_A x) = b\}) = \text{card } \{x \in A. f' (p_A x) = b\}$ **for** b
proof –
have *inj-on* $p_A \{x \in A. f' (p_A x) = b\}$

```

    by (metis (no-types, lifting) ⟨pA permutes A⟩ injD inj-onI permutes-inj)
    from this show ?thesis by (simp add: card-image)
qed
moreover have pA ‘ {x ∈ A. f' (pA x) = b} = {x ∈ A. f' x = b} for b
proof
  show pA ‘ {x ∈ A. f' (pA x) = b} ⊆ {x ∈ A. f' x = b}
    by (auto simp add: ⟨pA permutes A⟩ permutes-in-image)
  show {x ∈ A. f' x = b} ⊆ pA ‘ {x ∈ A. f' (pA x) = b}
    proof
      fix x
      assume x ∈ {x ∈ A. f' x = b}
      moreover have pA (inv pA x) = x
        using ⟨pA permutes A⟩ permutes-inverses(1) by fastforce
      moreover from ⟨x ∈ {x ∈ A. f' x = b}⟩ have inv pA x ∈ A
        by (simp add: ⟨pA permutes A⟩ permutes-in-image permutes-inv)
      ultimately show x ∈ pA ‘ {x ∈ A. f' (pA x) = b}
        by (auto intro: image-eqI[where x=inv pA x])
    qed
  qed
  ultimately show ?thesis by auto
qed
show (λX. inv pA ‘ X) ‘ {X ∈ (λb. {x ∈ A. f' x = b}) ‘ B - {∅}. card X
= n} ⊆ {X ∈ (λb. {x ∈ A. f' (pA x) = b}) ‘ B - {∅}. card X = n}
proof -
  have card (inv pA ‘ {x ∈ A. f' x = b}) = card {x ∈ A. f' x = b} for b
  proof -
    have inj-on (inv pA) {x ∈ A. f' x = b}
      by (metis (no-types, lifting) ⟨pA permutes A⟩ injD inj-onI permutes-surj
surj-imp-inj-inv)
    from this show ?thesis by (simp add: card-image)
  qed
  moreover have inv pA ‘ {x ∈ A. f' x = b} = {x ∈ A. f' (pA x) = b} for b
  proof
    show inv pA ‘ {x ∈ A. f' x = b} ⊆ {x ∈ A. f' (pA x) = b}
      using ⟨pA permutes A⟩
    by (auto simp add: permutes-in-image permutes-inv permutes-inverses(1))
    show {x ∈ A. f' (pA x) = b} ⊆ inv pA ‘ {x ∈ A. f' x = b}
      proof
        fix x
        assume x ∈ {x ∈ A. f' (pA x) = b}
        moreover have inv pA (pA x) = x
          by (meson ⟨pA permutes A⟩ permutes-inverses(2))
        moreover from ⟨x ∈ {x ∈ A. f' (pA x) = b}⟩ have pA x ∈ A
          by (simp add: ⟨pA permutes A⟩ permutes-in-image)
        ultimately show x ∈ inv pA ‘ {x ∈ A. f' x = b}
          by (auto intro: image-eqI[where x=pA x])
      qed
    qed
  qed
  ultimately show ?thesis by auto

```



```

    qed
  qed
  from this have card  $\{x' \in (\lambda b. \{x \in A. f' (p_A x) = b\}) \text{ ' } B - \{\{\}\}. \text{ card } x' = n\}$ 
    = card  $\{x' \in (\lambda b. \{x \in A. f' x = b\}) \text{ ' } B - \{\{\}\}. \text{ card } x' = n\}$ 
    by (rule bij-betw-same-card)
  from this show count (image-mset card (mset-set (( $\lambda b. \{x \in A. f' (p_A x) = b\}$ )
    '  $B - \{\{\}\}$ )))  $n =$ 
    count (image-mset card (mset-set (( $\lambda b. \{x \in A. f' x = b\}$ ) '  $B - \{\{\}\}$ )))  $n$ 
    using <finite B> by (simp add: count-image-mset')
  qed
  finally show image-mset card (mset-set (( $\lambda b. \{x \in A. f x = b\}$ ) '  $B - \{\{\}\}$ )) =
    image-mset card (mset-set (( $\lambda b. \{x \in A. f' x = b\}$ ) '  $B - \{\{\}\}$ )) .
  qed
end

```

4 Functions from A to B

```

theory Twelvefold-Way-Entry1
imports Preliminaries
begin

```

Note that the cardinality theorems of both structures, lists and finite functions, are already available. Hence, this development creates the bijection between those two structures and transfers the one cardinality theorem to the other structures and vice versa, although not strictly needed as both cardinality theorems were already available.

4.1 Definition of Bijections

```

definition sequence-of :: 'a set  $\Rightarrow$  (nat  $\Rightarrow$  'a)  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  'b list
where
  sequence-of A enum f = map ( $\lambda n. f (enum n)$ ) [0.. $\text{card } A$ ]

```

```

definition function-of :: 'a set  $\Rightarrow$  (nat  $\Rightarrow$  'a)  $\Rightarrow$  'b list  $\Rightarrow$  ('a  $\Rightarrow$  'b)
where
  function-of A enum xs = ( $\lambda a. \text{if } a \in A \text{ then } xs ! \text{inv-into } \{0.. $\text{length } xs\} \text{ enum } a \text{ else undefined}$ )$ 
```

4.2 Properties for Bijections

```

lemma nth-sequence-of:
  assumes  $i < \text{card } A$ 
  shows (sequence-of A enum f) !  $i = f (enum i)$ 
using assms unfolding sequence-of-def by auto

```

```

lemma nth-sequence-of-inv-into:
  assumes bij-betw enum  $\{0.. $\text{card } A\}$  A$ 
```

```

assumes  $a \in A$ 
shows  $(\text{sequence-of } A \text{ enum } f) ! (\text{inv-into } \{0..<\text{card } A\} \text{ enum } a) = f a$ 
proof -
  have  $\text{inv-into } \{0..<\text{card } A\} \text{ enum } a \in \{0..<\text{card } A\}$ 
  using  $\text{assms bij-betwE bij-betw-inv-into}$  by  $\text{blast}$ 
  from  $\text{this assms}$  show  $(\text{sequence-of } A \text{ enum } f) ! (\text{inv-into } \{0..<\text{card } A\} \text{ enum } a)$ 
   $= f a$ 
  unfolding  $\text{sequence-of-def}$  by  $(\text{simp add: bij-betw-inv-into-right})$ 
qed

```

```

lemma  $\text{set-sequence-of}$ :
  assumes  $\text{bij-betw enum } \{0..<\text{card } A\} A$ 
  assumes  $f \in A \rightarrow_E B$ 
  shows  $\text{set } (\text{sequence-of } A \text{ enum } f) \subseteq B$ 
using  $\text{PiE bij-betwE assms}$ 
unfolding  $\text{sequence-of-def}$  by  $\text{fastforce}$ 

```

```

lemma  $\text{length-sequence-of}$ :
  assumes  $\text{bij-betw enum } \{0..<\text{card } A\} A$ 
  assumes  $f \in A \rightarrow_E B$ 
  shows  $\text{length } (\text{sequence-of } A \text{ enum } f) = \text{card } A$ 
using  $\text{assms}$  unfolding  $\text{sequence-of-def}$  by  $\text{simp}$ 

```

```

lemma  $\text{function-of-enum}$ :
  assumes  $\text{bij-betw enum } \{0..<\text{card } A\} A$ 
  assumes  $\text{length } xs = \text{card } A$ 
  assumes  $i < \text{card } A$ 
  shows  $\text{function-of } A \text{ enum } xs (\text{enum } i) = xs ! i$ 
using  $\text{assms}$  unfolding  $\text{function-of-def}$ 
by  $(\text{auto simp add: bij-betw-inv-into-left bij-betwE})$ 

```

```

lemma  $\text{function-of-in-extensional-funcset}$ :
  assumes  $\text{bij-betw enum } \{0..<\text{card } A\} A$ 
  assumes  $\text{set } xs \subseteq B \text{ length } xs = \text{card } A$ 
  shows  $\text{function-of } A \text{ enum } xs \in A \rightarrow_E B$ 
proof
  fix  $x$ 
  assume  $x \in A$ 
  have  $\text{inv-into } \{0..<\text{length } xs\} \text{ enum } x \in \{0..<\text{length } xs\}$ 
  using  $\langle x \in A \rangle \text{ assms}(1, 3)$  by  $(\text{metis bij-betw-def inv-into-into})$ 
  from  $\text{this}$  have  $xs ! \text{inv-into } \{0..<\text{length } xs\} \text{ enum } x \in \text{set } xs$  by  $\text{simp}$ 
  from  $\text{this } \langle \text{set } xs \subseteq B \rangle$  show  $\text{function-of } A \text{ enum } xs x \in B$ 
  using  $\langle x \in A \rangle$  unfolding  $\text{function-of-def}$  by  $\text{auto}$ 
next
  fix  $x$ 
  assume  $x \notin A$ 
  from  $\text{this}$  show  $\text{function-of } A \text{ enum } xs x = \text{undefined}$ 
  unfolding  $\text{function-of-def}$  by  $\text{simp}$ 
qed

```

lemma *sequence-of-function-of*:
assumes *bij-betw* *enum* $\{0..<\text{card } A\}$ *A*
assumes *set* $xs \subseteq B$ *length* $xs = \text{card } A$
shows *sequence-of* *A enum* (*function-of* *A enum* xs) = xs
proof (*rule nth-equalityI*)
have *function-of* *A enum* $xs \in A \rightarrow_E B$
using *assms* **by** (*rule function-of-in-extensional-funcset*)
from this **show** *length* (*sequence-of* *A enum* (*function-of* *A enum* xs)) = *length* xs
using *assms*(1,3) **by** (*simp add: length-sequence-of*)
from this **show** $\bigwedge i. i < \text{length} (\text{sequence-of } A \text{ enum } (\text{function-of } A \text{ enum } xs))$
 $\Rightarrow \text{sequence-of } A \text{ enum } (\text{function-of } A \text{ enum } xs) ! i = xs ! i$
using *assms* **by** (*auto simp add: nth-sequence-of function-of-enum*)
qed

lemma *function-of-sequence-of*:
assumes *bij-betw* *enum* $\{0..<\text{card } A\}$ *A*
assumes $f \in A \rightarrow_E B$
shows *function-of* *A enum* (*sequence-of* *A enum* f) = f
proof
fix x
show *function-of* *A enum* (*sequence-of* *A enum* f) $x = f x$
using *assms* **unfolding** *function-of-def*
by (*auto simp add: length-sequence-of nth-sequence-of-inv-into*)
qed

4.3 Bijections

lemma *bij-betw-sequence-of*:
assumes *bij-betw* *enum* $\{0..<\text{card } A\}$ *A*
shows *bij-betw* (*sequence-of* *A enum*) ($A \rightarrow_E B$) $\{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\}$
proof (*rule bij-betw-byWitness*[**where** $f' = \text{function-of } A \text{ enum}$])
show $\forall f \in A \rightarrow_E B. \text{function-of } A \text{ enum } (\text{sequence-of } A \text{ enum } f) = f$
using *assms* **by** (*simp add: function-of-sequence-of*)
show $\forall xs \in \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\}. \text{sequence-of } A \text{ enum } (\text{function-of } A \text{ enum } xs) = xs$
using *assms* **by** (*auto simp add: sequence-of-function-of*)
show *sequence-of* *A enum* ' $(A \rightarrow_E B) \subseteq \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\}$
using *assms* *set-sequence-of*[*OF* *assms*] *length-sequence-of* **by** *auto*
show *function-of* *A enum* ' $\{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\} \subseteq A \rightarrow_E B$
using *assms* *function-of-in-extensional-funcset* **by** *blast*
qed

lemma *bij-betw-function-of*:
assumes *bij-betw* *enum* $\{0..<\text{card } A\}$ *A*
shows *bij-betw* (*function-of* *A enum*) $\{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\}$ ($A \rightarrow_E B$)

proof (rule *bij-betw-byWitness*[**where** $f' = \text{sequence-of } A \text{ enum}$])
show $\forall f \in A \rightarrow_E B. \text{function-of } A \text{ enum } (\text{sequence-of } A \text{ enum } f) = f$
using *assms* **by** (simp add: function-of-sequence-of)
show $\forall xs \in \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\}. \text{sequence-of } A \text{ enum } (\text{function-of } A \text{ enum } xs) = xs$
using *assms* **by** (auto simp add: sequence-of-function-of)
show $\text{sequence-of } A \text{ enum } \langle A \rightarrow_E B \rangle \subseteq \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\}$
using *assms* *set-sequence-of*[*OF assms*] *length-sequence-of* **by** auto
show $\text{function-of } A \text{ enum } \langle \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\} \rangle \subseteq A \rightarrow_E B$
using *assms* *function-of-in-extensional-funcset* **by** blast
qed

4.4 Cardinality

lemma

assumes *finite A*

shows $\text{card } (A \rightarrow_E B) = \text{card } B \wedge \text{card } A$

proof –

obtain *enum* **where** *bij-betw enum* $\{0..<\text{card } A\} A$

using $\langle \text{finite } A \rangle$ *ex-bij-betw-nat-finite* **by** blast

have *bij-betw* $(\text{sequence-of } A \text{ enum}) (A \rightarrow_E B) \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\}$

using $\langle \text{bij-betw enum } \{0..<\text{card } A\} A \rangle$ **by** (rule *bij-betw-sequence-of*)

from this **have** $\text{card } (A \rightarrow_E B) = \text{card } \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\}$

by (rule *bij-betw-same-card*)

also have $\text{card } \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\} = \text{card } B \wedge \text{card } A$

by (rule *card-lists-length-eq*)

finally show *?thesis* .

qed

lemma *card-sequences*:

assumes *finite A*

shows $\text{card } \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\} = \text{card } B \wedge \text{card } A$

proof –

obtain *enum* **where** *bij-betw enum* $\{0..<\text{card } A\} A$

using $\langle \text{finite } A \rangle$ *ex-bij-betw-nat-finite* **by** blast

have *bij-betw* $(\text{function-of } A \text{ enum}) \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\} (A \rightarrow_E B)$

using $\langle \text{bij-betw enum } \{0..<\text{card } A\} A \rangle$ **by** (rule *bij-betw-function-of*)

from this **have** $\text{card } \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A\} = \text{card } (A \rightarrow_E B)$

by (rule *bij-betw-same-card*)

also have $\text{card } (A \rightarrow_E B) = \text{card } B \wedge \text{card } A$

using $\langle \text{finite } A \rangle$ **by** (rule *card-extensional-funcset*)

finally show *?thesis* .

qed

lemma

shows $\text{card } \{xs. \text{set } xs \subseteq A \wedge \text{length } xs = n\} = \text{card } A \wedge n$

proof –

```

    have card {xs. set xs  $\subseteq$  A  $\wedge$  length xs = n} = card {xs. set xs  $\subseteq$  A  $\wedge$  length xs
= card {0.. $n$ }}
    by auto
    also have ... = card A  $\wedge$  card {0.. $n$ } by (subst card-sequences) auto
    also have ... = card A  $\wedge$  n by auto
    finally show ?thesis .
qed

end

```

5 Injections from A to B

```

theory Twelfold-Way-Entry2
imports Twelfold-Way-Entry1
begin

```

Note that the cardinality theorems of both structures, distinct lists and finite injective functions, are already available. Hence, this development creates the bijection between those two structures and transfers the one cardinality theorem to the other structures and vice versa, although not strictly needed as both cardinality theorems were already available.

5.1 Properties for Bijections

```

lemma inj-on-implies-distinct:
  assumes bij-betw enum {0.. $\text{card } A$ } A
  assumes  $f \in A \rightarrow_E B$ 
  assumes inj-on f A
  shows distinct (sequence-of A enum f)
proof -
  {
    fix i j
    assume bounds:  $i < \text{length } (\text{sequence-of } A \text{ enum } f) \wedge j < \text{length } (\text{sequence-of } A \text{ enum } f)$ 
    assume  $i \neq j$ 
    from bounds assms(1, 2) have bounds':  $i < \text{card } A \wedge j < \text{card } A$ 
    using length-sequence-of by fastforce+
    from this assms(1) have in-A:  $\text{enum } i \in A \wedge \text{enum } j \in A$ 
    using bij-betwE by fastforce+
    from  $\langle i \neq j \rangle$  bounds' assms(1) have enum  $i \neq \text{enum } j$ 
    by (metis bij-betw-inv-into-left lessThan-iff atLeast0LessThan)
    from this have  $f (\text{enum } i) \neq f (\text{enum } j)$ 
    using assms(3) in-A inj-onD by fastforce
    from this bounds' have sequence-of A enum  $f ! i \neq \text{sequence-of } A \text{ enum } f ! j$ 
    by (simp add: nth-sequence-of)
  }
from this show ?thesis
by (auto simp add: distinct-conv-nth)

```

qed

lemma *distinct-implies-inj-on*:

assumes *bij-betw* *enum* $\{0..<\text{card } A\}$ *A*
assumes *length xs = card A*
assumes *distinct xs*
shows *inj-on* (*function-of A enum xs*) *A*
proof (*rule inj-onI*)
let $?idx\text{-of} = \lambda x. \text{inv-into } \{0..<\text{length } xs\} \text{ enum } x$
fix *x y*
assume $x \in A \ y \in A \ \text{function-of } A \ \text{enum } xs \ x = \text{function-of } A \ \text{enum } xs \ y$
from *this* **have** $xs ! ?idx\text{-of } x = xs ! ?idx\text{-of } y$
unfolding *function-of-def* **by** *simp*
have $?idx\text{-of } x = ?idx\text{-of } y$
proof –
have $?idx\text{-of } x < \text{length } xs$
using $\langle x \in A \rangle \text{ assms}(1,2)$
by (*metis atLeast0LessThan bij-betw-imp-surj-on inv-into-into lessThan-iff*)
moreover **have** $?idx\text{-of } y < \text{length } xs$
using $\langle y \in A \rangle \text{ assms}(1,2)$
by (*metis atLeast0LessThan bij-betw-imp-surj-on inv-into-into lessThan-iff*)
moreover **note** $\langle xs ! ?idx\text{-of } x = xs ! ?idx\text{-of } y \rangle \langle \text{distinct } xs \rangle$
ultimately **show** *?thesis*
by (*auto dest: nth-eq-iff-index-eq[where i=?idx-of x and j=?idx-of y]*)
qed
from *this* $\langle \text{bij-betw} \ - \ - \rightarrow \rangle$ **show** $x = y$
by (*metis* $\langle x \in A \rangle \langle y \in A \rangle \langle \text{length } xs = \text{card } A \rangle \text{bij-betw-inv-into-right}$)
qed

lemma *image-sequence-of-inj*:

assumes *bij-betw* *enum* $\{0..<\text{card } A\}$ *A*
shows *sequence-of A enum* $\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} \subseteq \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}$
proof
fix *xs*
assume $xs \in \text{sequence-of } A \ \text{enum } \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$
from *this* **obtain** *f* **where** $xs = \text{sequence-of } A \ \text{enum } f$ **and** $f: f \in A \rightarrow_E B$
inj-on f A **by** *auto*
moreover **from** $xs \ f \ \langle \text{bij-betw} \ - \ - \rightarrow \rangle$ **have** $\text{set } xs \subseteq B$
using *set-sequence-of subsetCE* **by** *blast*
moreover **from** $xs \ f \ \langle \text{bij-betw} \ - \ - \rightarrow \rangle$ **have** $\text{length } xs = \text{card } A$
using *length-sequence-of* **by** *auto*
moreover **from** $xs \ f \ \langle \text{bij-betw} \ - \ - \rightarrow \rangle$ **have** *distinct xs*
using *inj-on-implies-distinct* **by** *simp*
ultimately **show** $xs \in \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}$ **by** *auto*
qed

lemma *image-function-of-distinct*:

assumes *bij-betw enum* $\{0..<\text{card } A\} \ A$
shows *function-of A enum* ‘ $\{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}$
 $\subseteq \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$
proof
fix f
assume $f: f \in \text{function-of } A \text{ enum } \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}$
from $f \text{ assms}$ **have** $f \in A \rightarrow_E B$
using *function-of-in-extensional-funcset* **by** *blast*
moreover from $f \text{ assms}$ **have** *inj-on* $f \ A$
by (*auto simp add: assms distinct-implies-inj-on*)
ultimately show $f \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$ **by** *auto*
qed

5.2 Bijections

lemma *bij-betw-sequence-of*:
assumes *bij-betw enum* $\{0..<\text{card } A\} \ A$
shows *bij-betw (sequence-of A enum)* $\{f. f \in A \rightarrow_E B \wedge \text{inj-on } f \ A\} \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}$
proof (*rule bij-betw-byWitness[where f'=function-of A enum]*)
show $\forall f \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}. \text{function-of } A \text{ enum } (\text{sequence-of } A \text{ enum } f) = f$
using *assms* **by** (*auto simp add: function-of-sequence-of*)
show $\forall xs \in \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}. \text{sequence-of } A \text{ enum } (\text{function-of } A \text{ enum } xs) = xs$
using *assms* **by** (*auto simp add: sequence-of-function-of*)
show *sequence-of A enum* ‘ $\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} \subseteq \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}$
using *assms* **by** (*simp add: image-sequence-of-inj*)
show *function-of A enum* ‘ $\{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\} \subseteq \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$
using *assms* **by** (*simp add: image-function-of-distinct*)
qed

lemma *bij-betw-function-of*:
assumes *bij-betw enum* $\{0..<\text{card } A\} \ A$
shows *bij-betw (function-of A enum)* $\{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\} \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$
proof (*rule bij-betw-byWitness[where f'=sequence-of A enum]*)
show $\forall f \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}. \text{function-of } A \text{ enum } (\text{sequence-of } A \text{ enum } f) = f$
using *assms* **by** (*auto simp add: function-of-sequence-of*)
show $\forall xs \in \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}. \text{sequence-of } A \text{ enum } (\text{function-of } A \text{ enum } xs) = xs$
using *assms* **by** (*auto simp add: sequence-of-function-of*)
show *sequence-of A enum* ‘ $\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} \subseteq \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}$
using *assms* **by** (*simp add: image-sequence-of-inj*)

show *function-of* A *enum* ‘ $\{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}$
 $\subseteq \{f \in A \rightarrow_E B. \text{inj-on } f A\}$
using *assms* **by** (*simp add: image-function-of-distinct*)
qed

5.3 Cardinality

lemma

assumes *finite* A *finite* B $\text{card } A \leq \text{card } B$
shows $\text{card } \{f \in A \rightarrow_E B. \text{inj-on } f A\} = \prod \{\text{card } B - \text{card } A + 1.. \text{card } B\}$
proof –
obtain *enum* **where** *bij-betw* *enum* $\{0..<\text{card } A\}$ A
using $\langle \text{finite } A \rangle$ *ex-bij-betw-nat-finite* **by** *blast*
have *bij-betw* (*sequence-of* A *enum*) $\{f \in A \rightarrow_E B. \text{inj-on } f A\}$ $\{xs. \text{set } xs \subseteq B$
 $\wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\}$
using $\langle \text{bij-betw } \text{enum } \{0..<\text{card } A\} A \rangle$ **by** (*rule bij-betw-sequence-of*)
from *this* **have** $\text{card } \{f \in A \rightarrow_E B. \text{inj-on } f A\} = \text{card } \{xs. \text{set } xs \subseteq B \wedge \text{length}$
 $xs = \text{card } A \wedge \text{distinct } xs\}$
by (*rule bij-betw-same-card*)
also have $\text{card } \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\} = \text{card } \{xs.$
 $\text{length } xs = \text{card } A \wedge \text{distinct } xs \wedge \text{set } xs \subseteq B\}$
by *meson*
also have $\text{card } \{xs. \text{length } xs = \text{card } A \wedge \text{distinct } xs \wedge \text{set } xs \subseteq B\} = \prod \{\text{card}$
 $B - \text{card } A + 1.. \text{card } B\}$
using $\langle \text{finite } B \rangle$ $\langle \text{card } A \leq \text{card } B \rangle$ **by** (*rule List.card-lists-distinct-length-eq*)
finally show *?thesis* .
qed

lemma *card-sequences:*

assumes *finite* A *finite* B $\text{card } A \leq \text{card } B$
shows $\text{card } \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\} = \text{fact } (\text{card } B)$
 $\text{div fact } (\text{card } B - \text{card } A)$
proof –
obtain *enum* **where** *bij-betw* *enum* $\{0..<\text{card } A\}$ A
using $\langle \text{finite } A \rangle$ *ex-bij-betw-nat-finite* **by** *blast*
have *bij-betw* (*function-of* A *enum*) $\{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge$
 $\text{distinct } xs\}$ $\{f \in A \rightarrow_E B. \text{inj-on } f A\}$
using $\langle \text{bij-betw } \text{enum } \{0..<\text{card } A\} A \rangle$ **by** (*rule bij-betw-function-of*)
from *this* **have** $\text{card } \{xs. \text{set } xs \subseteq B \wedge \text{length } xs = \text{card } A \wedge \text{distinct } xs\} = \text{card}$
 $\{f \in A \rightarrow_E B. \text{inj-on } f A\}$
by (*rule bij-betw-same-card*)
also have $\text{card } \{f \in A \rightarrow_E B. \text{inj-on } f A\} = \text{fact } (\text{card } B) \text{ div fact } (\text{card } B -$
 $\text{card } A)$
using $\langle \text{finite } A \rangle$ $\langle \text{finite } B \rangle$ $\langle \text{card } A \leq \text{card } B \rangle$ **by** (*rule card-extensional-funcset-inj-on*)
finally show *?thesis* .
qed

end

6 Functions from A to B, up to a Permutation of A

```
theory Twelvefold-Way-Entry4
imports Equiv-Relations-on-Functions
begin
```

6.1 Definition of Bijections

```
definition msubset-of :: 'a set  $\Rightarrow$  ('a  $\Rightarrow$  'b) set  $\Rightarrow$  'b multiset
where
  msubset-of A F = univ ( $\lambda$ f. image-mset f (mset-set A)) F
```

```
definition functions-of :: 'a set  $\Rightarrow$  'b multiset  $\Rightarrow$  ('a  $\Rightarrow$  'b) set
where
  functions-of A B = {f  $\in$  A  $\rightarrow_E$  set-mset B. image-mset f (mset-set A) = B}
```

6.2 Properties for Bijections

```
lemma msubset-of:
  assumes F  $\in$  (A  $\rightarrow_E$  B) // domain-permutation A B
  shows size (msubset-of A F) = card A
  and set-mset (msubset-of A F)  $\subseteq$  B
proof -
  from  $\langle F \in (A \rightarrow_E B) // \text{domain-permutation } A B \rangle$  obtain f where f  $\in$  A  $\rightarrow_E$  B
  and F-eq: F = domain-permutation A B “ {f} using quotientE by blast
  have msubset-of A F = univ ( $\lambda$ f. image-mset f (mset-set A)) F
  unfolding msubset-of-def ..
  also have ... = univ ( $\lambda$ f. image-mset f (mset-set A)) (domain-permutation A B “ {f})
  unfolding F-eq ..
  also have ... = image-mset f (mset-set A)
  using equiv-domain-permutation image-mset-respects-domain-permutation  $\langle f \in$ 
A  $\rightarrow_E$  B  $\rangle$ 
  by (subst univ-commute') auto
  finally have msubset-of-eq: msubset-of A F = image-mset f (mset-set A) .
  show size (msubset-of A F) = card A
  proof -
    have size (msubset-of A F) = size (image-mset f (mset-set A))
    unfolding msubset-of-eq ..
    also have ... = card A
    by (cases  $\langle \text{finite } A \rangle$ ) auto
    finally show ?thesis .
  qed
  show set-mset (msubset-of A F)  $\subseteq$  B
  proof -
    have set-mset (msubset-of A F) = set-mset (image-mset f (mset-set A))
    unfolding msubset-of-eq ..
```

```

    also have ...  $\subseteq$  B
    using  $\langle f \in A \rightarrow_E B \rangle$  by (cases finite A) auto
    finally show ?thesis .
qed
qed

lemma functions-of:
  assumes finite A
  assumes set-mset  $M \subseteq B$ 
  assumes size  $M = \text{card } A$ 
  shows functions-of A  $M \in (A \rightarrow_E B)$  // domain-permutation A B
proof -
  obtain f where  $f \in A \rightarrow_E \text{set-mset } M$  and  $\text{image-mset } f (\text{mset-set } A) = M$ 
  using obtain-function-on-ext-funcset  $\langle \text{finite } A \rangle \langle \text{size } M = \text{card } A \rangle$  by blast
  from  $\langle f \in A \rightarrow_E \text{set-mset } M \rangle$  have  $f \in A \rightarrow_E B$ 
  using  $\langle \text{set-mset } M \subseteq B \rangle$  PiE-iff subset-eq by blast
  have functions-of A  $M = (\text{domain-permutation } A B)$  “ {f}
  proof
    show functions-of A  $M \subseteq \text{domain-permutation } A B$  “ {f}
    proof
      fix f'
      assume  $f' \in \text{functions-of } A M$ 
      from this have  $M = \text{image-mset } f' (\text{mset-set } A)$  and  $f' \in A \rightarrow_E f' ' A$ 
      using  $\langle \text{finite } A \rangle$  unfolding functions-of-def by auto
      from this assms(1, 2) have  $f' \in A \rightarrow_E B$ 
      by (simp add: PiE-iff image-subset-iff)
      obtain p where p permutes A  $\wedge (\forall x \in A. f x = f' (p x))$ 
      using  $\langle \text{finite } A \rangle \langle \text{image-mset } f (\text{mset-set } A) = M \rangle \langle M = \text{image-mset } f' (\text{mset-set } A) \rangle$ 
      image-mset-eq-implyes-permutes by blast
      from this show  $f' \in \text{domain-permutation } A B$  “ {f}
      using  $\langle f \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle$ 
      unfolding domain-permutation-def by auto
    qed
  next
    show domain-permutation A B “ {f}  $\subseteq \text{functions-of } A M$ 
    proof
      fix f'
      assume  $f' \in \text{domain-permutation } A B$  “ {f}
      from this have  $(f, f') \in \text{domain-permutation } A B$  by auto
      from this  $\langle \text{image-mset } f (\text{mset-set } A) = M \rangle$  have  $\text{image-mset } f' (\text{mset-set } A) = M$ 
      using congruentD[OF image-mset-respects-domain-permutation] by metis
      moreover from this  $\langle (f, f') \in \text{domain-permutation } A B \rangle$  have  $f' \in A \rightarrow_E \text{set-mset } M$ 
      using  $\langle \text{finite } A \rangle$  unfolding domain-permutation-def by auto
      ultimately show  $f' \in \text{functions-of } A M$ 
      unfolding functions-of-def by auto
    qed
  qed

```

```

qed
from this  $\langle f \in A \rightarrow_E B \rangle$  show ?thesis by (auto intro: quotientI)
qed

lemma functions-of-msubset-of:
  assumes finite A
  assumes  $F \in (A \rightarrow_E B)$  // domain-permutation A B
  shows functions-of A (msubset-of A F) = F
proof -
  from  $\langle F \in (A \rightarrow_E B) // \text{domain-permutation A B} \rangle$  obtain f where  $f \in A \rightarrow_E B$ 
  and F-eq:  $F = \text{domain-permutation A B} \text{ `` } \{f\}$  using quotientE by blast
  have msubset-of A F = univ  $(\lambda f. \text{image-mset } f (\text{mset-set A}))$  F
  unfolding msubset-of-def ..
  also have  $\dots = \text{univ } (\lambda f. \text{image-mset } f (\text{mset-set A})) (\text{domain-permutation A B} \text{ `` } \{f\})$ 
  unfolding F-eq ..
  also have  $\dots = \text{image-mset } f (\text{mset-set A})$ 
  using equiv-domain-permutation image-mset-respects-domain-permutation  $\langle f \in A \rightarrow_E B \rangle$ 
  by (subst univ-commute') auto
  finally have msubset-of-eq:  $\text{msubset-of A F} = \text{image-mset } f (\text{mset-set A})$  .
  show ?thesis
proof
  show functions-of A (msubset-of A F)  $\subseteq$  F
proof
  fix f'
  assume  $f' \in \text{functions-of A (msubset-of A F)}$ 
  from this have  $f': f' \in A \rightarrow_E f' \text{ `` set-mset (mset-set A)}$ 
  image-mset  $f' (\text{mset-set A}) = \text{image-mset } f (\text{mset-set A})$ 
  unfolding functions-of-def by (auto simp add: msubset-of-eq)
  from  $\langle f \in A \rightarrow_E B \rangle$  have  $f' \text{ `` } A \subseteq B$  by auto
  note  $\langle f \in A \rightarrow_E B \rangle$ 
  moreover from  $f'(1) \langle \text{finite A} \rangle \langle f' \text{ `` } A \subseteq B \rangle$  have  $f' \in A \rightarrow_E B$  by auto
  moreover obtain p where p permutes A  $\wedge (\forall x \in A. f x = f' (p x))$ 
  using  $\langle \text{finite A} \rangle \langle \text{image-mset } f' (\text{mset-set A}) = \text{image-mset } f (\text{mset-set A}) \rangle$ 
  by (metis image-mset-eq-implies-permutes)
  ultimately show  $f' \in F$ 
  unfolding F-eq domain-permutation-def by auto
qed
next
show  $F \subseteq \text{functions-of A (msubset-of A F)}$ 
proof
  fix f'
  assume  $f' \in F$ 
  from this have  $f' \in A \rightarrow_E B$ 
  unfolding F-eq domain-permutation-def by auto
  from  $\langle f' \in F \rangle$  obtain p where p permutes A  $\wedge (\forall x \in A. f x = f' (p x))$ 
  unfolding F-eq domain-permutation-def by auto

```

```

from this have eq: image-mset f' (mset-set A) = image-mset f (mset-set A)
  using permutes-implies-image-mset-eq by blast
moreover have  $f' \in A \rightarrow_E \text{set-mset } (\text{image-mset } f \text{ (mset-set } A))$ 
  using  $\langle \text{finite } A \rangle \langle f' \in A \rightarrow_E B \rangle \text{eq[symmetric]}$  by auto
ultimately show  $f' \in \text{functions-of } A \text{ (msubset-of } A \text{ } F)$ 
  unfolding functions-of-def msubset-of-eq by auto
qed
qed
qed

lemma msubset-of-functions-of:
  assumes  $\text{set-mset } M \subseteq B \text{ size } M = \text{card } A \text{ finite } A$ 
  shows  $\text{msubset-of } A \text{ (functions-of } A \text{ } M) = M$ 
proof –
  from assms have  $\text{functions-of } A \text{ } M \in (A \rightarrow_E B) // \text{domain-permutation } A \text{ } B$ 
  using functions-of by fastforce
  from this obtain f where  $f \in A \rightarrow_E B$  and  $\text{functions-of } A \text{ } M = \text{domain-permutation}$ 
   $A \text{ } B \text{ “}\{f\}$ 
  by (rule quotientE)
  from this have  $f \in \text{functions-of } A \text{ } M$ 
  using equiv-domain-permutation equiv-class-self by fastforce
  have  $\text{msubset-of } A \text{ (functions-of } A \text{ } M) = \text{univ } (\lambda f. \text{image-mset } f \text{ (mset-set } A))$ 
  (functions-of } A \text{ } M)
  unfolding msubset-of-def ..
  also have  $\dots = \text{univ } (\lambda f. \text{image-mset } f \text{ (mset-set } A)) \text{ (domain-permutation } A \text{ } B$ 
   $\text{“}\{f\})$ 
  unfolding  $\langle \text{functions-of } A \text{ } M = \text{domain-permutation } A \text{ } B \text{ “}\{f\} \rangle ..$ 
  also have  $\dots = \text{image-mset } f \text{ (mset-set } A)$ 
  using equiv-domain-permutation image-mset-respects-domain-permutation  $\langle f \in$ 
 $A \rightarrow_E B \rangle$ 
  by (subst univ-commute') auto
  also have  $\text{image-mset } f \text{ (mset-set } A) = M$ 
  using  $\langle f \in \text{functions-of } A \text{ } M \rangle$  unfolding functions-of-def by simp
  finally show ?thesis .
qed

```

6.3 Bijections

```

lemma bij-betw-msubset-of:
  assumes finite A
  shows  $\text{bij-betw } (\text{msubset-of } A) ((A \rightarrow_E B) // \text{domain-permutation } A \text{ } B) \{M. \text{set-mset } M \subseteq B \wedge \text{size } M = \text{card } A\}$ 
proof (rule bij-betw-byWitness[where f'= $\lambda M. \text{functions-of } A \text{ } M$ ])
  show  $\forall F \in (A \rightarrow_E B) // \text{domain-permutation } A \text{ } B. \text{functions-of } A \text{ (msubset-of } A \text{ } F) = F$ 
  using  $\langle \text{finite } A \rangle$  by (auto simp add: functions-of-msubset-of)
  show  $\forall M \in \{M. \text{set-mset } M \subseteq B \wedge \text{size } M = \text{card } A\}. \text{msubset-of } A \text{ (functions-of } A \text{ } M) = M$ 
  using  $\langle \text{finite } A \rangle$  by (auto simp add: msubset-of-functions-of)

```

```

show msubset-of A ‘  $((A \rightarrow_E B) // \text{domain-permutation } A \ B) \subseteq \{M. \text{set-mset } M \subseteq B \wedge \text{size } M = \text{card } A\}$ 
using msubset-of by blast
show functions-of A ‘  $\{M. \text{set-mset } M \subseteq B \wedge \text{size } M = \text{card } A\} \subseteq (A \rightarrow_E B)$ 
// domain-permutation A B
using functions-of  $\langle \text{finite } A \rangle$  by blast
qed

```

6.4 Cardinality

lemma

```

assumes finite A finite B
shows  $\text{card } ((A \rightarrow_E B) // \text{domain-permutation } A \ B) = \text{card } B + \text{card } A - 1$ 
choose card A
proof –
have bij-betw (msubset-of A)  $((A \rightarrow_E B) // \text{domain-permutation } A \ B)$   $\{M. \text{set-mset } M \subseteq B \wedge \text{size } M = \text{card } A\}$ 
using  $\langle \text{finite } A \rangle$  by (rule bij-betw-msubset-of)
from this have  $\text{card } ((A \rightarrow_E B) // \text{domain-permutation } A \ B) = \text{card } \{M. \text{set-mset } M \subseteq B \wedge \text{size } M = \text{card } A\}$ 
by (rule bij-betw-same-card)
also have  $\text{card } \{M. \text{set-mset } M \subseteq B \wedge \text{size } M = \text{card } A\} = \text{card } B + \text{card } A - 1$ 
choose card A
using  $\langle \text{finite } B \rangle$  by (rule card-multisets)
finally show ?thesis .
qed

```

end

7 Injections from A to B up to a Permutation of A

theory *Twelvefold-Way-Entry5*

imports

Equiv-Relations-on-Functions

begin

7.1 Definition of Bijections

definition *subset-of* :: $'a \text{ set} \Rightarrow ('a \Rightarrow 'b) \text{ set} \Rightarrow 'b \text{ set}$

where

$\text{subset-of } A \ F = \text{univ } (\lambda f. f \text{ ‘ } A) \ F$

definition *functions-of* :: $'a \text{ set} \Rightarrow 'b \text{ set} \Rightarrow ('a \Rightarrow 'b) \text{ set}$

where

$\text{functions-of } A \ B = \{f \in A \rightarrow_E B. f \text{ ‘ } A = B\}$

7.2 Properties for Bijections

lemma *functions-of-eq*:

assumes *finite A*

assumes $f \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$

shows *functions-of A (f ‘ A) = domain-permutation A B “ {f}*

proof

have *bij*: *bij-betw f A (f ‘ A)*

using *assms* **by** (*simp add: bij-betw-imageI*)

show *functions-of A (f ‘ A) \subseteq domain-permutation A B “ {f}*

proof

fix *f'*

assume $f' \in \text{functions-of } A \ (f' \text{ ‘ } A)$

from *this* **have** $f' \in A \rightarrow_E f' \text{ ‘ } A$ **and** $f' \text{ ‘ } A = f' \text{ ‘ } A$

unfolding *functions-of-def* **by** *auto*

from *this* *assms* **have** $f' \in A \rightarrow_E B$ **and** *inj-on f A*

using *PiE-mem* **by** *fastforce+*

moreover **have** $\exists p. p \text{ permutes } A \wedge (\forall x \in A. f \ x = f' \ (p \ x))$

proof

let $?p = \lambda x. \text{if } x \in A \text{ then } \text{inv-into } A \ f' \ (f \ x) \text{ else } x$

show $?p \text{ permutes } A \wedge (\forall x \in A. f \ x = f' \ (?p \ x))$

proof

show $?p \text{ permutes } A$

proof (*rule bij-imp-permutes*)

show *bij-betw ?p A A*

proof (*rule bij-betw-imageI*)

show *inj-on ?p A*

proof (*rule inj-onI*)

fix *a a'*

assume $a \in A \ a' \in A \ ?p \ a = ?p \ a'$

from *this* **have** $\text{inv-into } A \ f' \ (f \ a) = \text{inv-into } A \ f' \ (f \ a')$ **by** *auto*

from *this* $\langle a \in A \rangle \langle a' \in A \rangle \langle f' \text{ ‘ } A = f' \text{ ‘ } A \rangle$ **have** $f \ a = f \ a'$

using *inv-into-injective* **by** *fastforce*

from *this* $\langle a \in A \rangle \langle a' \in A \rangle$ **show** $a = a'$

by (*metis bij bij-betw-inv-into-left*)

qed

next

show $?p \text{ ‘ } A = A$

proof

show $?p \text{ ‘ } A \subseteq A$

using $\langle f' \text{ ‘ } A = f' \text{ ‘ } A \rangle$ **by** (*simp add: image-subsetI inv-into-into*)

next

show $A \subseteq ?p \text{ ‘ } A$

proof

fix *a*

assume $a \in A$

have *inj-on f' A*

using $\langle \text{finite } A \rangle \langle f' \text{ ‘ } A = f' \text{ ‘ } A \rangle \langle \text{inj-on } f \ A \rangle$

by (*simp add: card-image eq-card-imp-inj-on*)

from $\langle a \in A \rangle \langle f' \text{ ‘ } A = f' \text{ ‘ } A \rangle$ **have** $\text{inv-into } A \ f \ (f' \ a) \in A$

```

      by (metis image-eqI inv-into-into)
    moreover have  $a = \text{inv-into } A \ f' \ (f \ (\text{inv-into } A \ f \ (f' \ a)))$ 
      using  $\langle a \in A \rangle \langle f' \text{ ' } A = f \text{ ' } A \rangle \langle \text{inj-on } f' \ A \rangle$ 
      by (metis f-inv-into-f image-eqI inv-into-f-f)
    ultimately show  $a \in ?p \text{ ' } A$  by auto
  qed
qed
qed
next
  fix  $x$ 
  assume  $x \notin A$ 
  from this show  $?p \ x = x$  by simp
qed
next
  from  $\langle f' \text{ ' } A = f \text{ ' } A \rangle$  show  $\forall x \in A. f \ x = f' \ (?p \ x)$ 
    by (simp add: f-inv-into-f)
  qed
qed
moreover have  $f \in A \rightarrow_E B$  using assms by auto
ultimately show  $f' \in \text{domain-permutation } A \ B \text{ '' } \{f\}$ 
  unfolding domain-permutation-def by auto
qed
next
show  $\text{domain-permutation } A \ B \text{ '' } \{f\} \subseteq \text{functions-of } A \ (f \text{ ' } A)$ 
proof
  fix  $f'$ 
  assume  $f' \in \text{domain-permutation } A \ B \text{ '' } \{f\}$ 
  from this obtain  $p$  where  $p$ :  $p$  permutes  $A \ \forall x \in A. f \ x = f' \ (p \ x)$ 
    and  $f \in A \rightarrow_E B \ f' \in A \rightarrow_E B$ 
    unfolding domain-permutation-def by auto
  have  $f' \text{ ' } A = f \text{ ' } A$ 
  proof
    show  $f' \text{ ' } A \subseteq f \text{ ' } A$ 
    proof
      fix  $x$ 
      assume  $x \in f' \text{ ' } A$ 
      from this obtain  $x'$  where  $x = f' \ x'$  and  $x' \in A$  ..
      from this have  $x = f \ (\text{inv } p \ x')$ 
      using  $p$  by (metis (mono-tags, lifting) permutes-in-image permutes-inverses(1))
      moreover have  $\text{inv } p \ x' \in A$ 
        using  $p \ \langle x' \in A \rangle$  by (simp add: permutes-in-image permutes-inv)
      ultimately show  $x \in f \text{ ' } A$  ..
    qed
  next
    show  $f \text{ ' } A \subseteq f' \text{ ' } A$ 
      using  $p$  permutes-in-image by fastforce
  qed
moreover from this  $\langle f' \in A \rightarrow_E B \rangle$  have  $f' \in A \rightarrow_E f \text{ ' } A$  by auto
ultimately show  $f' \in \text{functions-of } A \ (f \text{ ' } A)$ 

```

unfolding functions-of-def by auto
 qed
 qed

lemma subset-of:
 assumes $F \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$ // domain-permutation $A \ B$
 shows $\text{subset-of } A \ F \subseteq B$ and $\text{card } (\text{subset-of } A \ F) = \text{card } A$
proof –
 from *assms* obtain f where $F\text{-eq}$: $F = (\text{domain-permutation } A \ B) \ \text{“}\{f\}$
 and $f: f \in A \rightarrow_E B \text{inj-on } f \ A$
 using *mem-Collect-eq quotientE* by *force*
 from *this* have $\text{subset-of } A \ (\text{domain-permutation } A \ B \ \text{“}\{f\}) = f \ ^\text{‘} A$
 using *equiv-domain-permutation image-respects-domain-permutation*
 unfolding *subset-of-def* by (*intro univ-commute'*) *auto*
 from *this* $f \ F\text{-eq}$ show $\text{subset-of } A \ F \subseteq B$ and $\text{card } (\text{subset-of } A \ F) = \text{card } A$
 by (*auto simp add: card-image*)
 qed

lemma functions-of:
 assumes *finite* A *finite* B $X \subseteq B$ $\text{card } X = \text{card } A$
 shows $\text{functions-of } A \ X \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$ // domain-permutation $A \ B$
proof –
 from *assms* obtain f where $f: f \in A \rightarrow_E X \wedge \text{bij-betw } f \ A \ X$
 using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ by (*metis finite-same-card-bij-on-ext-funcset finite-subset*)
 from *this* have $X = f \ ^\text{‘} A$ by (*simp add: bij-betw-def*)
 from $f \ \langle X \subseteq B \rangle$ have $f \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$
 by (*auto simp add: bij-betw-imp-inj-on*)
 have $\text{functions-of } A \ X = \text{domain-permutation } A \ B \ \text{“}\{f\}$
 using $\langle \text{finite } A \rangle \langle X = f \ ^\text{‘} A \rangle \langle f \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\} \rangle$
 by (*simp add: functions-of-eq*)
 from *this* show $\text{functions-of } A \ X \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$ // domain-permutation $A \ B$
 using $\langle f \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\} \rangle$ by (*auto intro: quotientI*)
 qed

lemma subset-of-functions-of:
 assumes *finite* A *finite* X $\text{card } A = \text{card } X$
 shows $\text{subset-of } A \ (\text{functions-of } A \ X) = X$
proof –
 from *assms* obtain f where $f \in A \rightarrow_E X$ and $\text{bij-betw } f \ A \ X$
 using *finite-same-card-bij-on-ext-funcset* by *blast*
 from *this* have $\text{subset-of: } \text{subset-of } A \ (\text{domain-permutation } A \ X \ \text{“}\{f\}) = f \ ^\text{‘} A$
 using *equiv-domain-permutation image-respects-domain-permutation*
 unfolding *subset-of-def* by (*intro univ-commute'*) *auto*
 from $\langle \text{bij-betw } f \ A \ X \rangle$ have $\text{inj-on } f \ A$ and $f \ ^\text{‘} A = X$
 by (*auto simp add: bij-betw-def*)
 have $\text{subset-of } A \ (\text{functions-of } A \ X) = \text{subset-of } A \ (\text{functions-of } A \ (f \ ^\text{‘} A))$


```

    using ⟨f ‘ A = X⟩ by simp
  also have ... = subset-of A (domain-permutation A X “ {f})
    using ⟨finite A⟩ ⟨inj-on f A⟩ ⟨f ∈ A →E X⟩ by (auto simp add: functions-of-eq)
  also have ... = f ‘ A
    using ⟨inj-on f A⟩ ⟨f ∈ A →E X⟩ by (simp add: subset-of)
  also have ... = X
    using ⟨f ‘ A = X⟩ by simp
  finally show ?thesis .
qed

```

```

lemma functions-of-subset-of:
  assumes finite A
  assumes F ∈ {f ∈ A →E B. inj-on f A} // domain-permutation A B
  shows functions-of A (subset-of A F) = F
using assms(2) proof (rule quotientE)
  fix f
  assume f: f ∈ {f ∈ A →E B. inj-on f A}
  and F-eq: F = domain-permutation A B “ {f}
  from this have subset-of A (domain-permutation A B “ {f}) = f ‘ A
    using equiv-domain-permutation image-respects-domain-permutation
    unfolding subset-of-def by (intro univ-commute′) auto
  from this f F-eq ⟨finite A⟩ show functions-of A (subset-of A F) = F
    by (simp add: functions-of-eq)
qed

```

7.3 Bijections

```

lemma bij-betw-subset-of:
  assumes finite A finite B
  shows bij-betw (subset-of A) ({f ∈ A →E B. inj-on f A} // domain-permutation
A B) {X. X ⊆ B ∧ card X = card A}
proof (rule bij-betw-byWitness[where f′=functions-of A])
  show ∀ F ∈ {f ∈ A →E B. inj-on f A} // domain-permutation A B. functions-of
A (subset-of A F) = F
    using ⟨finite A⟩ functions-of-subset-of by auto
  show ∀ X ∈ {X. X ⊆ B ∧ card X = card A}. subset-of A (functions-of A X) = X
    using subset-of-functions-of ⟨finite A⟩ ⟨finite B⟩
    by (metis (mono-tags) finite-subset mem-Collect-eq)
  show subset-of A “ ({f ∈ A →E B. inj-on f A} // domain-permutation A B) ⊆
{X. X ⊆ B ∧ card X = card A}
    using subset-of by fastforce
  show functions-of A “ {X. X ⊆ B ∧ card X = card A} ⊆ {f ∈ A →E B. inj-on
f A} // domain-permutation A B
    using ⟨finite A⟩ ⟨finite B⟩ functions-of by auto
qed

```

```

lemma bij-betw-functions-of:
  assumes finite A finite B
  shows bij-betw (functions-of A) {X. X ⊆ B ∧ card X = card A} ({f ∈ A →E

```

$B. \text{inj-on } f A \} // \text{domain-permutation } A B)$
proof (rule *bij-betw-byWitness*[**where** $f' = \text{subset-of } A$])
show $\forall F \in \{f \in A \rightarrow_E B. \text{inj-on } f A \} // \text{domain-permutation } A B. \text{functions-of } A (\text{subset-of } A F) = F$
using $\langle \text{finite } A \rangle \text{functions-of-subset-of}$ **by** *auto*
show $\forall X \in \{X. X \subseteq B \wedge \text{card } X = \text{card } A\}. \text{subset-of } A (\text{functions-of } A X) = X$
using *subset-of-functions-of* $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$
by (*metis* (*mono-tags*) *finite-subset mem-Collect-eq*)
show $\text{subset-of } A ' (\{f \in A \rightarrow_E B. \text{inj-on } f A \} // \text{domain-permutation } A B) \subseteq \{X. X \subseteq B \wedge \text{card } X = \text{card } A\}$
using *subset-of* **by** *fastforce*
show $\text{functions-of } A ' \{X. X \subseteq B \wedge \text{card } X = \text{card } A\} \subseteq \{f \in A \rightarrow_E B. \text{inj-on } f A \} // \text{domain-permutation } A B$
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle \text{functions-of}$ **by** *auto*
qed

lemma *bij-betw-mset-set*:

shows *bij-betw mset-set* $\{A. \text{finite } A\} \{M. \forall x. \text{count } M x \leq 1\}$
proof (rule *bij-betw-byWitness*[**where** $f' = \text{set-mset}$])
show $\forall A \in \{A. \text{finite } A\}. \text{set-mset } (\text{mset-set } A) = A$ **by** *auto*
show $\forall M \in \{M. \forall x. \text{count } M x \leq 1\}. \text{mset-set } (\text{set-mset } M) = M$
by (*auto simp add: mset-set-set-mset*)
show $\text{mset-set } ' \{A. \text{finite } A\} \subseteq \{M. \forall x. \text{count } M x \leq 1\}$
using *nat-le-linear* **by** *fastforce*
show $\text{set-mset } ' \{M. \forall x. \text{count } M x \leq 1\} \subseteq \{A. \text{finite } A\}$ **by** *auto*
qed

lemma *bij-betw-mset-set-card*:

assumes *finite A*
shows *bij-betw mset-set* $\{X. X \subseteq A \wedge \text{card } X = k\} \{M. M \subseteq \# \text{mset-set } A \wedge \text{size } M = k\}$
proof (rule *bij-betw-byWitness*[**where** $f' = \text{set-mset}$])
show $\forall X \in \{X. X \subseteq A \wedge \text{card } X = k\}. \text{set-mset } (\text{mset-set } X) = X$
using $\langle \text{finite } A \rangle \text{rev-finite-subset[of } A]$ **by** *auto*
show $\forall M \in \{M. M \subseteq \# \text{mset-set } A \wedge \text{size } M = k\}. \text{mset-set } (\text{set-mset } M) = M$
by (*auto simp add: mset-set-set-mset*)
show $\text{mset-set } ' \{X. X \subseteq A \wedge \text{card } X = k\} \subseteq \{M. M \subseteq \# \text{mset-set } A \wedge \text{size } M = k\}$
using $\langle \text{finite } A \rangle \text{rev-finite-subset[of } A]$
by (*auto simp add: mset-set-subseteq-mset-set*)
show $\text{set-mset } ' \{M. M \subseteq \# \text{mset-set } A \wedge \text{size } M = k\} \subseteq \{X. X \subseteq A \wedge \text{card } X = k\}$
using *assms mset-subset-eqD card-set-mset* **by** *fastforce*
qed

lemma *bij-betw-mset-set-card'*:

assumes *finite A*
shows *bij-betw mset-set* $\{X. X \subseteq A \wedge \text{card } X = k\} \{M. \text{set-mset } M \subseteq A \wedge \text{size } M = k \wedge (\forall x. \text{count } M x \leq 1)\}$

proof (rule *bij-betw-byWitness*[**where** $f' = \text{set-mset}$])
show $\forall X \in \{X. X \subseteq A \wedge \text{card } X = k\}. \text{set-mset } (\text{mset-set } X) = X$
using $\langle \text{finite } A \rangle \text{ rev-finite-subset[of } A]$ **by** *auto*
show $\forall M \in \{M. \text{set-mset } M \subseteq A \wedge \text{size } M = k \wedge (\forall x. \text{count } M \ x \leq 1)\}. \text{mset-set } (\text{set-mset } M) = M$
by (auto simp add: *mset-set-set-mset'*)
show $\text{mset-set } \{X. X \subseteq A \wedge \text{card } X = k\} \subseteq \{M. \text{set-mset } M \subseteq A \wedge \text{size } M = k \wedge (\forall x. \text{count } M \ x \leq 1)\}$
using $\langle \text{finite } A \rangle \text{ rev-finite-subset[of } A]$ **by** (auto simp add: *count-mset-set-leq'*)
show $\text{set-mset } \{M. \text{set-mset } M \subseteq A \wedge \text{size } M = k \wedge (\forall x. \text{count } M \ x \leq 1)\} \subseteq \{X. X \subseteq A \wedge \text{card } X = k\}$
by (auto simp add: *card-set-mset'*)
qed

7.4 Cardinality

lemma *card-injective-functions-domain-permutation:*

assumes *finite A*
shows $\text{card } (\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-permutation } A \ B) = \text{card } B$
choose card A
proof –
have *bij-betw (subset-of A) ({f ∈ A →_E B. inj-on f A} // domain-permutation A B) {X. X ⊆ B ∧ card X = card A}*
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ **by** (rule *bij-betw-subset-of*)
from this have $\text{card } (\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-permutation } A \ B) = \text{card } \{X. X \subseteq B \wedge \text{card } X = \text{card } A\}$
by (rule *bij-betw-same-card*)
also have $\text{card } \{X. X \subseteq B \wedge \text{card } X = \text{card } A\} = \text{card } B$ *choose card A*
using $\langle \text{finite } B \rangle$ **by** (rule *n-subsets*)
finally show *?thesis* .
qed

lemma *card-multiset-only-sets:*

assumes *finite A*
shows $\text{card } \{M. M \subseteq\# \text{mset-set } A \wedge \text{size } M = k\} = \text{card } A$ *choose k*
proof –
have *bij-betw mset-set {X. X ⊆ A ∧ card X = k} {M. M ⊆# mset-set A ∧ size M = k}*
using $\langle \text{finite } A \rangle$ **by** (rule *bij-betw-mset-set-card*)
from this have $\text{card } \{M. M \subseteq\# \text{mset-set } A \wedge \text{size } M = k\} = \text{card } \{X. X \subseteq A \wedge \text{card } X = k\}$
by (simp add: *bij-betw-same-card*)
also have $\text{card } \{X. X \subseteq A \wedge \text{card } X = k\} = \text{card } A$ *choose k*
using $\langle \text{finite } A \rangle$ **by** (rule *n-subsets*)
finally show *?thesis* .
qed

lemma *card-multiset-only-sets':*

assumes *finite A*

```

shows  $\text{card } \{M. \text{ set-mset } M \subseteq A \wedge \text{size } M = k \wedge (\forall x. \text{count } M x \leq 1)\} = \text{card}$ 
 $A \text{ choose } k$ 
proof –
  from  $\langle \text{finite } A \rangle$  have  $\{M. \text{ set-mset } M \subseteq A \wedge \text{size } M = k \wedge (\forall x. \text{count } M x \leq$ 
 $1)\} =$ 
 $\{M. M \subseteq\# \text{ mset-set } A \wedge \text{size } M = k\}$ 
  using msubset-mset-set-iff by auto
  from this  $\langle \text{finite } A \rangle$  card-multiset-only-sets show ?thesis by simp
qed

end

```

8 Surjections from A to B up to a Permutation on A

```

theory Twelvefold-Way-Entry6
imports Twelvefold-Way-Entry4
begin

```

8.1 Properties for Bijections

```

lemma set-mset-eq-implies-surj-on:
  assumes finite A
  assumes  $\text{size } M = \text{card } A \text{ set-mset } M = B$ 
  assumes  $f \in \text{functions-of } A \ M$ 
  shows  $f ' A = B$ 
proof –
  from  $\langle f \in \text{functions-of } A \ M \rangle$  have  $\text{image-mset } f \ (\text{mset-set } A) = M$ 
  unfolding functions-of-def by auto
  from  $\langle \text{image-mset } f \ (\text{mset-set } A) = M \rangle$  show  $f ' A = B$ 
  using  $\langle \text{set-mset } M = B \rangle \langle \text{finite } A \rangle$  finite-set-mset-mset-set set-image-mset by
force
qed

```

```

lemma surj-on-implies-set-mset-eq:
  assumes finite A
  assumes  $F \in (A \rightarrow_E B) \ /\ / \text{domain-permutation } A \ B$ 
  assumes  $\text{univ } (\lambda f. f ' A = B) \ F$ 
  shows  $\text{set-mset } (\text{msubset-of } A \ F) = B$ 
proof –
  from  $\langle F \in (A \rightarrow_E B) \ /\ / \text{domain-permutation } A \ B \rangle$  obtain  $f$  where  $f \in A \rightarrow_E B$ 
  and  $F\text{-eq}: F = \text{domain-permutation } A \ B \text{ “}\{f\}$  using quotientE by blast
  have  $\text{msubset-of } A \ F = \text{univ } (\lambda f. \text{image-mset } f \ (\text{mset-set } A)) \ F$ 
  unfolding msubset-of-def ..
  also have  $\dots = \text{univ } (\lambda f. \text{image-mset } f \ (\text{mset-set } A)) \ (\text{domain-permutation } A \ B$ 
 $\text{“}\{f\})$ 
  unfolding  $F\text{-eq}$  ..

```

```

also have ... = image-mset f (mset-set A)
using equiv-domain-permutation image-mset-respects-domain-permutation  $\langle f \in$ 
 $A \rightarrow_E B \rangle$ 
by (subst univ-commute') auto
finally have eq: msubset-of A F = image-mset f (mset-set A) .
from iffD1[OF univ-commute', OF equiv-domain-permutation, OF surjective-respects-domain-permutation,
 $OF \langle f \in A \rightarrow_E B \rangle$ ]
 $\langle univ (\lambda f. f ' A = B) F \rangle$  have  $f ' A = B$  by (simp add: F-eq)
have set-mset (image-mset f (mset-set A)) = B
proof
show set-mset (image-mset f (mset-set A))  $\subseteq B$ 
using  $\langle finite A \rangle \langle f ' A = B \rangle$  by auto
next
show  $B \subseteq set-mset (image-mset f (mset-set A))$ 
using  $\langle finite A \rangle$  by (simp add:  $\langle f ' A = B \rangle$ [symmetric] in-image-mset)
qed
from this show set-mset (msubset-of A F) = B
unfolding eq .
qed

```

lemma *functions-of-is-surj-on*:

```

assumes finite A
assumes size M = card A set-mset M = B
shows univ  $(\lambda f. f ' A = B)$  (functions-of A M)
proof –
have functions-of A M  $\in (A \rightarrow_E B)$  // domain-permutation A B
using functions-of  $\langle finite A \rangle \langle size M = card A \rangle \langle set-mset M = B \rangle$  by fastforce
from this obtain f where eq-f: functions-of A M = domain-permutation A B
“  $\{f\}$  and  $f \in A \rightarrow_E B$ 
using quotientE by blast
from eq-f have  $f \in functions-of A M$ 
using  $\langle f \in A \rightarrow_E B \rangle$  equiv-domain-permutation equiv-class-self by fastforce
have  $f ' A = B$ 
using  $\langle f \in functions-of A M \rangle$  assms set-mset-eq-implies-surj-on by fastforce
from this show ?thesis
unfolding eq-f using equiv-domain-permutation surjective-respects-domain-permutation
 $\langle f \in A \rightarrow_E B \rangle$ 
by (subst univ-commute') assumption+
qed

```

8.2 Bijections

lemma *bij-betw-msubset-of*:

```

assumes finite A
shows bij-betw (msubset-of A) ( $\{f \in A \rightarrow_E B. f ' A = B\}$  // domain-permutation
 $A B$ )
 $\{M. set-mset M = B \wedge size M = card A\}$ 
(is bij-betw - ?FSet ?MSet)
proof (rule bij-betw-byWitness[where  $f'=\lambda M. functions-of A M$ ])

```

```

have quotient-eq: ?FSet = {F ∈ ((A →E B) // domain-permutation A B). univ
(λf. f ‘ A = B) F}
using equiv-domain-permutation[of A B] surjective-respects-domain-permutation[of
A B]
by (simp only: univ-preserves-predicate)
show ∀f ∈ ?FSet. functions-of A (msubset-of A f) = f
using ⟨finite A⟩ by (auto simp only: quotient-eq functions-of-msubset-of)
show ∀M ∈ ?MSet. msubset-of A (functions-of A M) = M
using ⟨finite A⟩ msubset-of-functions-of by blast
show msubset-of A ‘ ?FSet ⊆ ?MSet
using ⟨finite A⟩ by (auto simp add: quotient-eq surj-on-implies-set-mset-eq
msubset-of)
show functions-of A ‘ ?MSet ⊆ ?FSet
using ⟨finite A⟩ by (auto simp add: quotient-eq intro: functions-of func-
tions-of-is-surj-on)
qed

```

8.3 Cardinality

```

lemma card-surjective-functions-domain-permutation:
assumes finite A finite B
assumes card B ≤ card A
shows card ({f ∈ A →E B. f ‘ A = B} // domain-permutation A B) = (card A
− 1) choose (card A − card B)
proof −
let ?FSet = {f ∈ A →E B. f ‘ A = B} // domain-permutation A B
and ?MSet = {M. set-mset M = B ∧ size M = card A}
have bij-betw (msubset-of A) ?FSet ?MSet
using ⟨finite A⟩ by (rule bij-betw-msubset-of)
from this have card ?FSet = card ?MSet
by (rule bij-betw-same-card)
also have card ?MSet = (card A − 1) choose (card A − card B)
using ⟨finite B⟩ ⟨card B ≤ card A⟩ by (rule card-multisets-covering-set)
finally show ?thesis .
qed

end

```

9 Functions from A to B up to a Permutation on B

```

theory Twelfefold-Way-Entry7
imports Equiv-Relations-on-Functions
begin

```

9.1 Definition of Bijections

```

definition partitions-of :: 'a set ⇒ 'b set ⇒ ('a ⇒ 'b) set ⇒ 'a set set
where

```

$\text{partitions-of } A \ B \ F = \text{univ } (\lambda f. (\lambda b. \{x \in A. f \ x = b\}) \text{ ' } B - \{\{\}\}) \ F$

definition $\text{functions-of} :: 'a \ \text{set} \ \text{set} \Rightarrow 'a \ \text{set} \Rightarrow 'b \ \text{set} \Rightarrow ('a \Rightarrow 'b) \ \text{set}$
where

$\text{functions-of } P \ A \ B = \{f \in A \rightarrow_E B. (\lambda b. \{x \in A. f \ x = b\}) \text{ ' } B - \{\{\}\} = P\}$

9.2 Properties for Bijections

lemma partitions-of :

assumes $\text{finite } B$

assumes $F \in (A \rightarrow_E B) \ // \ \text{range-permutation } A \ B$

shows $\text{card } (\text{partitions-of } A \ B \ F) \leq \text{card } B$

and $\text{partition-on } A \ (\text{partitions-of } A \ B \ F)$

proof –

from $\langle F \in (A \rightarrow_E B) \ // \ \text{range-permutation } A \ B \rangle$ **obtain** f **where** $f \in A \rightarrow_E B$

and $F\text{-eq}$: $F = \text{range-permutation } A \ B \text{ " } \{f\} \text{ using quotientE by blast}$

have $\text{partitions-of } A \ B \ F = \text{univ } (\lambda f. (\lambda b. \{x \in A. f \ x = b\}) \text{ ' } B - \{\{\}\}) \ F$

unfolding $\text{partitions-of-def} \ ..$

also have $\dots = \text{univ } (\lambda f. (\lambda b. \{x \in A. f \ x = b\}) \text{ ' } B - \{\{\}\}) \ (\text{range-permutation } A \ B \text{ " } \{f\})$

unfolding $F\text{-eq} \ ..$

also have $\dots = (\lambda b. \{x \in A. f \ x = b\}) \text{ ' } B - \{\{\}\}$

using $\text{equiv-range-permutation domain-partitions-respects-range-permutation } \langle f \in A \rightarrow_E B \rangle$

by $(\text{subst univ-commute'}) \ \text{auto}$

finally have partitions-of-eq : $\text{partitions-of } A \ B \ F = (\lambda b. \{x \in A. f \ x = b\}) \text{ ' } B - \{\{\}\} \ .$

show $\text{card } (\text{partitions-of } A \ B \ F) \leq \text{card } B$

proof –

have $\text{card } (\text{partitions-of } A \ B \ F) = \text{card } ((\lambda b. \{x \in A. f \ x = b\}) \text{ ' } B - \{\{\}\})$

unfolding $\text{partitions-of-eq} \ ..$

also have $\dots \leq \text{card } ((\lambda b. \{x \in A. f \ x = b\}) \text{ ' } B)$

using $\langle \text{finite } B \rangle$ **by** $(\text{auto intro: card-mono})$

also have $\dots \leq \text{card } B$

using $\langle \text{finite } B \rangle$ **by** $(\text{rule card-image-le})$

finally show $?thesis \ .$

qed

show $\text{partition-on } A \ (\text{partitions-of } A \ B \ F)$

proof –

have $\text{partition-on } A \ ((\lambda b. \{x \in A. f \ x = b\}) \text{ ' } B - \{\{\}\})$

using $\langle f \in A \rightarrow_E B \rangle$ **by** $(\text{auto intro!: partition-onI})$

from this show $?thesis$

unfolding $\text{partitions-of-eq} \ .$

qed

qed

lemma functions-of :

assumes $\text{finite } A \ \text{finite } B$

```

    assumes partition-on  $A$   $P$ 
    assumes  $\text{card } P \leq \text{card } B$ 
    shows  $\text{functions-of } P \ A \ B \in (A \rightarrow_E B) \ // \ \text{range-permutation } A \ B$ 
  proof -
    obtain  $f$  where  $f \in A \rightarrow_E B$  and  $r1: (\lambda b. \{x \in A. f \ x = b\}) \text{ ` } B - \{\{\}\} = P$ 
    using obtain-function-with-partition[OF  $\langle \text{finite } A \rangle \langle \text{finite } B \rangle \langle \text{partition-on } A \ P \rangle$ 
     $\langle \text{card } P \leq \text{card } B \rangle$ ]
    by blast
    have  $\text{functions-of } P \ A \ B = \text{range-permutation } A \ B \text{ `` } \{f\}$ 
  proof
    show  $\text{functions-of } P \ A \ B \subseteq \text{range-permutation } A \ B \text{ `` } \{f\}$ 
  proof
    fix  $f'$ 
    assume  $f' \in \text{functions-of } P \ A \ B$ 
    from this have  $f' \in A \rightarrow_E B$  and  $r2: (\lambda b. \{x \in A. f' \ x = b\}) \text{ ` } B - \{\{\}\}$ 
    =  $P$ 
    unfolding functions-of-def by auto
    from  $r1 \ r2$ 
    obtain  $p$  where  $p$  permutes  $B \wedge (\forall x \in A. f \ x = p \ (f' \ x))$ 
    using partitions-eq-implies-permutes[OF  $\langle f \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle$ 
     $\langle \text{finite } B \rangle$  by metis
    from this show  $f' \in \text{range-permutation } A \ B \text{ `` } \{f\}$ 
    using  $\langle f \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle$ 
    unfolding range-permutation-def by auto
  qed
next
show  $\text{range-permutation } A \ B \text{ `` } \{f\} \subseteq \text{functions-of } P \ A \ B$ 
proof
  fix  $f'$ 
  assume  $f' \in \text{range-permutation } A \ B \text{ `` } \{f\}$ 
  from this have  $(f, f') \in \text{range-permutation } A \ B$  by auto
  from this have  $f' \in A \rightarrow_E B$ 
  unfolding range-permutation-def by auto
  from  $\langle (f, f') \in \text{range-permutation } A \ B \rangle$  have
     $(\lambda b. \{x \in A. f \ x = b\}) \text{ ` } B - \{\{\}\} = (\lambda b. \{x \in A. f' \ x = b\}) \text{ ` } B - \{\{\}\}$ 
  using congruentD[OF domain-partitions-respects-range-permutation] by blast
  from  $\langle f' \in A \rightarrow_E B \rangle$  this  $r1$  show  $f' \in \text{functions-of } P \ A \ B$ 
  unfolding functions-of-def by auto
qed
qed
from this  $\langle f \in A \rightarrow_E B \rangle$  show ?thesis by (auto intro: quotientI)
qed

lemma functions-of-partitions-of:
  assumes finite  $B$ 
  assumes  $F \in (A \rightarrow_E B) \ // \ \text{range-permutation } A \ B$ 
  shows  $\text{functions-of } (\text{partitions-of } A \ B \ F) \ A \ B = F$ 
proof -
  from  $\langle F \in (A \rightarrow_E B) \ // \ \text{range-permutation } A \ B \rangle$  obtain  $f$  where  $f \in A \rightarrow_E B$ 

```



```

B
  and F-eq: F = range-permutation A B “ {f} using quotientE by blast
  have partitions-of-eq: partitions-of A B F = (λb. {x ∈ A. f x = b}) ‘ B - {[]}
    unfolding partitions-of-def F-eq
    using equiv-range-permutation domain-partitions-respects-range-permutation
  ⟨f ∈ A →E B⟩
    by (subst univ-commute′) auto
  show ?thesis
  proof
    show functions-of (partitions-of A B F) A B ⊆ F
    proof
      fix f′
      assume f′: f′ ∈ functions-of (partitions-of A B F) A B
      from this have (λb. {x ∈ A. f x = b}) ‘ B - {[]} = (λb. {x ∈ A. f′ x = b})
    ‘ B - {[]}
      unfolding functions-of-def by (auto simp add: partitions-of-eq)
      note ⟨f ∈ A →E B⟩
      moreover from f′ have f′ ∈ A →E B
      unfolding functions-of-def by auto
      moreover obtain p where p permutes B ∧ (∀ x ∈ A. f x = p (f′ x))
      using partitions-eq-implies-permutes[OF ⟨f ∈ A →E B⟩ ⟨f′ ∈ A →E B⟩
    ⟨finite B⟩
      ⟨(λb. {x ∈ A. f x = b}) ‘ B - {[]} = (λb. {x ∈ A. f′ x = b}) ‘ B - {[]}⟩
      by metis
      ultimately show f′ ∈ F
      unfolding F-eq range-permutation-def by auto
    qed
  next
    show F ⊆ functions-of (partitions-of A B F) A B
    proof
      fix f′
      assume f′ ∈ F
      from this have f′ ∈ A →E B
      unfolding F-eq range-permutation-def by auto
      from ⟨f′ ∈ F⟩ obtain p where p permutes B ∧ ∀ x ∈ A. f x = p (f′ x)
      unfolding F-eq range-permutation-def by auto
      have eq: (λb. {x ∈ A. f′ x = b}) ‘ B - {[]} = (λb. {x ∈ A. f x = b}) ‘ B -
    {}
      proof -
        have (λb. {x ∈ A. f′ x = b}) ‘ B - {[]} = (λb. {x ∈ A. p (f′ x) = b}) ‘ B
      - {}
        using permutes-implies-inv-image-on-eq[OF ⟨p permutes B⟩, of A f′] by
    simp
      also have ... = (λb. {x ∈ A. f x = b}) ‘ B - {}
      using ⟨∀ x ∈ A. f x = p (f′ x)⟩ by auto
      finally show ?thesis .
    qed
    from this ⟨f′ ∈ A →E B⟩ show f′ ∈ functions-of (partitions-of A B F) A B
    unfolding functions-of-def partitions-of-eq by auto

```

qed
 qed
 qed

lemma *partitions-of-functions-of*:

assumes *finite A finite B*
assumes *partition-on A P*
assumes *card P ≤ card B*
shows *partitions-of A B (functions-of P A B) = P*
proof –
have *functions-of P A B ∈ (A →_E B) // range-permutation A B*
using *⟨finite A⟩ ⟨finite B⟩ ⟨partition-on A P⟩ ⟨card P ≤ card B⟩* **by** (*rule functions-of*)
from this obtain f where *f ∈ A →_E B* **and** *functions-of-eq: functions-of P A B = range-permutation A B “{f}”*
using *quotientE* **by** *metis*
from *functions-of-eq ⟨f ∈ A →_E B⟩* **have** *f ∈ functions-of P A B*
using *equiv-range-permutation equiv-class-self* **by** *fastforce*
have *partitions-of A B (functions-of P A B) = univ (λf. (λb. {x ∈ A. f x = b}) ‘B – {f}’) (functions-of P A B)*
unfolding *partitions-of-def ..*
also have *... = univ (λf. (λb. {x ∈ A. f x = b}) ‘B – {f}’) (range-permutation A B “{f}”)*
unfolding *⟨functions-of P A B = range-permutation A B “{f}” ..*
also have *... = (λb. {x ∈ A. f x = b}) ‘B – {f}’*
using *equiv-range-permutation domain-partitions-respects-range-permutation ⟨f ∈ A →_E B⟩*
by (*subst univ-commute’*) *auto*
also have *(λb. {x ∈ A. f x = b}) ‘B – {f}’ = P*
using *⟨f ∈ functions-of P A B⟩* **unfolding** *functions-of-def* **by** *simp*
finally show *?thesis .*
 qed

9.3 Bijections

lemma *bij-betw-partitions-of*:

assumes *finite A finite B*
shows *bij-betw (partitions-of A B) ((A →_E B) // range-permutation A B) {P. partition-on A P ∧ card P ≤ card B}*
proof (*rule bij-betw-byWitness[where f'=λP. functions-of P A B]*)
show *∀ F ∈ (A →_E B) // range-permutation A B. functions-of (partitions-of A B F) A B = F*
using *⟨finite B⟩* **by** (*simp add: functions-of-partitions-of*)
show *∀ P ∈ {P. partition-on A P ∧ card P ≤ card B}. partitions-of A B (functions-of P A B) = P*
using *⟨finite A⟩ ⟨finite B⟩* **by** (*auto simp add: partitions-of-functions-of*)
show *partitions-of A B ‘((A →_E B) // range-permutation A B) ⊆ {P. partition-on A P ∧ card P ≤ card B}*
using *⟨finite B⟩ partitions-of* **by** *auto*

show $(\lambda P. \text{functions-of } P \ A \ B) \subseteq \{P. \text{partition-on } A \ P \wedge \text{card } P \leq \text{card } B\} \subseteq$
 $(A \rightarrow_E B) \ // \ \text{range-permutation } A \ B$
using $\text{functions-of } \langle \text{finite } A \rangle \ \langle \text{finite } B \rangle$ **by** *auto*
qed

9.4 Cardinality

lemma

assumes $\text{finite } A \ \text{finite } B$
shows $\text{card } ((A \rightarrow_E B) \ // \ \text{range-permutation } A \ B) = (\sum j \leq \text{card } B. \text{Stirling } (\text{card } A) \ j)$
proof –
have $\text{bij-betw } (\text{partitions-of } A \ B) ((A \rightarrow_E B) \ // \ \text{range-permutation } A \ B) \ \{P. \text{partition-on } A \ P \wedge \text{card } P \leq \text{card } B\}$
using $\langle \text{finite } A \rangle \ \langle \text{finite } B \rangle$ **by** $(\text{rule } \text{bij-betw-partitions-of})$
from this have $\text{card } ((A \rightarrow_E B) \ // \ \text{range-permutation } A \ B) = \text{card } \{P. \text{partition-on } A \ P \wedge \text{card } P \leq \text{card } B\}$
by $(\text{rule } \text{bij-betw-same-card})$
also have $\text{card } \{P. \text{partition-on } A \ P \wedge \text{card } P \leq \text{card } B\} = (\sum j \leq \text{card } B. \text{Stirling } (\text{card } A) \ j)$
using $\langle \text{finite } A \rangle$ **by** $(\text{rule } \text{card-partition-on-at-most-size})$
finally show *?thesis* .
qed

end

10 Injections from A to B up to a Permutation on B

theory *Twelvefold-Way-Entry8*
imports *Twelvefold-Way-Entry7*
begin

10.1 Properties for Bijections

lemma *inj-on-implies-partitions-of*:

assumes $F \in (A \rightarrow_E B) \ // \ \text{range-permutation } A \ B$
assumes $\text{univ } (\lambda f. \text{inj-on } f \ A) \ F$
shows $\forall X \in \text{partitions-of } A \ B \ F. \text{card } X = 1$
proof –
from $\langle F \in (A \rightarrow_E B) \ // \ \text{range-permutation } A \ B \rangle$ **obtain** f **where** $f \in A \rightarrow_E B$
and $F\text{-eq: } F = \text{range-permutation } A \ B \ \{f\}$ **using** *quotientE* **by** *blast*
from this $\langle \text{univ } (\lambda f. \text{inj-on } f \ A) \ F \rangle$ **have** $\text{inj-on } f \ A$
using *univ-commute'[OF equiv-range-permutation inj-on-respects-range-permutation*
 $\langle f \in A \rightarrow_E B \rangle]$ **by** *simp*
have $\forall X \in (\lambda b. \{x \in A. f \ x = b\}) \ \{B - \{\{\}\}. \text{card } X = 1$
proof
fix X

```

    assume  $X \in (\lambda b. \{x \in A. f x = b\}) \text{ ' } B - \{\{\}\}$ 
    from this obtain  $x$  where  $X = \{xa \in A. f xa = f x\}$   $x \in A$  by auto
    from this have  $X = \{x\}$ 
      using  $\langle \text{inj-on } f \ A \rangle$  by (auto dest!: inj-onD)
    from this show  $\text{card } X = 1$  by simp
  qed
  from this show ?thesis
    unfolding partitions-of-def F-eq
    using equiv-range-permutation domain-partitions-respects-range-permutation  $\langle f \in A \rightarrow_E B \rangle$ 
    by (subst univ-commute') assumption+
  qed

```

lemma unique-part-eq-singleton:

```

  assumes partition-on  $A \ P$ 
  assumes  $\forall X \in P. \text{card } X = 1$ 
  assumes  $x \in A$ 
  shows  $(THE X. x \in X \wedge X \in P) = \{x\}$ 
proof -
  have  $(THE X. x \in X \wedge X \in P) \in P$ 
    using  $\langle \text{partition-on } A \ P \rangle \langle x \in A \rangle$  by (simp add: partition-on-the-part-mem)
  from this have  $\text{card } (THE X. x \in X \wedge X \in P) = 1$ 
    using  $\langle \forall X \in P. \text{card } X = 1 \rangle$  by auto
  moreover have  $x \in (THE X. x \in X \wedge X \in P)$ 
    using  $\langle \text{partition-on } A \ P \rangle \langle x \in A \rangle$  by (simp add: partition-on-in-the-unique-part)
  ultimately show ?thesis
    by (metis card-1-singletonE singleton-iff)
  qed

```

lemma functions-of-is-inj-on:

```

  assumes finite  $A$  finite  $B$  partition-on  $A \ P$   $\text{card } P \leq \text{card } B$ 
  assumes  $\forall X \in P. \text{card } X = 1$ 
  shows univ  $(\lambda f. \text{inj-on } f \ A)$  (functions-of  $P \ A \ B$ )
proof -
  have functions-of  $P \ A \ B \in (A \rightarrow_E B) // \text{range-permutation } A \ B$ 
    using functions-of  $\langle \text{finite } A \rangle \langle \text{finite } B \rangle \langle \text{partition-on } A \ P \rangle \langle \text{card } P \leq \text{card } B \rangle$ 
  by blast
  from this obtain  $f$  where eq-f: functions-of  $P \ A \ B = \text{range-permutation } A \ B$ 
  “  $\{f\}$  and  $f \in A \rightarrow_E B$ 
    using quotientE by blast
  from eq-f have  $f \in \text{functions-of } P \ A \ B$ 
    using  $\langle f \in A \rightarrow_E B \rangle$  equiv-range-permutation equiv-class-self by fastforce
  from this have eq:  $(\lambda b. \{x \in A. f x = b\}) \text{ ' } B - \{\{\}\} = P$ 
    unfolding functions-of-def by auto
  have inj-on  $f \ A$ 
  proof (rule inj-onI)
    fix  $x \ y$ 
    assume  $x \in A \ y \in A \ f x = f y$ 
    from  $\langle x \in A \rangle$  have  $x \in \{x' \in A. f x' = f x\}$  by auto

```

moreover from $\langle y \in A \rangle \langle f x = f y \rangle$ **have** $y \in \{x' \in A. f x' = f x\}$ **by** *auto*
moreover have $\text{card } \{x' \in A. f x' = f x\} = 1$
proof –
from $\langle x \in A \rangle \langle f \in A \rightarrow_E B \rangle$ **have** $f x \in B$ **by** *auto*
from this $\langle x \in A \rangle$ **have** $\{x' \in A. f x' = f x\} \in (\lambda b. \{x \in A. f x = b\}) \text{ ` } B -$
 $\{\{\}\}$ **by** *auto*
from this $\langle \forall X \in P. \text{card } X = 1 \rangle$ **eq show** *?thesis* **by** *auto*
qed
ultimately show $x = y$ **by** $(\text{metis card-1-singletonE singletonD})$
qed
from this show *?thesis*
unfolding *eq-f* **using** *equiv-range-permutation inj-on-respects-range-permutation*
 $\langle f \in A \rightarrow_E B \rangle$
by $(\text{subst univ-commute'})$ **assumption+**
qed

10.2 Bijections

lemma *bij-betw-partitions-of*:
assumes *finite A finite B*
shows *bij-betw (partitions-of A B) ($\{f \in A \rightarrow_E B. \text{inj-on } f A\}$ // range-permutation A B) $\{P. \text{partition-on } A P \wedge \text{card } P \leq \text{card } B \wedge (\forall X \in P. \text{card } X = 1)\}$*
proof $(\text{rule bij-betw-byWitness[where } f' = \lambda P. \text{functions-of } P A B])$
have *quotient-eq: $\{f \in A \rightarrow_E B. \text{inj-on } f A\}$ // range-permutation A B = $\{F \in ((A \rightarrow_E B) // \text{range-permutation A B}). \text{univ } (\lambda f. \text{inj-on } f A) F\}$*
by $(\text{simp add: equiv-range-permutation inj-on-respects-range-permutation univ-preserves-predicate})$
show $\forall F \in \{f \in A \rightarrow_E B. \text{inj-on } f A\} // \text{range-permutation A B}. \text{functions-of (partitions-of A B F) A B} = F$
using $\langle \text{finite B} \rangle$ **by** $(\text{simp add: quotient-eq functions-of-partitions-of})$
show $\forall P \in \{P. \text{partition-on } A P \wedge \text{card } P \leq \text{card } B \wedge (\forall X \in P. \text{card } X = 1)\}.$
 $\text{partitions-of A B (functions-of } P A B) = P$
using $\langle \text{finite A} \rangle \langle \text{finite B} \rangle$ **by** $(\text{simp add: partitions-of-functions-of})$
show *partitions-of A B $\text{' } (\{f \in A \rightarrow_E B. \text{inj-on } f A\} // \text{range-permutation A B})$*
 $\subseteq \{P. \text{partition-on } A P \wedge \text{card } P \leq \text{card } B \wedge (\forall X \in P. \text{card } X = 1)\}$
using $\langle \text{finite B} \rangle$ *quotient-eq partitions-of inj-on-implies-partitions-of* **by** *fastforce*
show $(\lambda P. \text{functions-of } P A B) \text{ ` } \{P. \text{partition-on } A P \wedge \text{card } P \leq \text{card } B \wedge (\forall X \in P. \text{card } X = 1)\} \subseteq \{f \in A \rightarrow_E B. \text{inj-on } f A\} // \text{range-permutation A B}$
using $\langle \text{finite A} \rangle \langle \text{finite B} \rangle$ **by** $(\text{auto simp add: quotient-eq intro: functions-of functions-of-is-inj-on})$
qed

10.3 Cardinality

lemma *card-injective-functions-range-permutation*:
assumes *finite A finite B*
shows $\text{card } (\{f \in A \rightarrow_E B. \text{inj-on } f A\} // \text{range-permutation A B}) = \text{iverson}(\text{card } A \leq \text{card } B)$
proof –
obtain *enum* **where** *bij-betw enum $\{0..<\text{card } A\}$ A*
using $\langle \text{finite A} \rangle$ *ex-bij-betw-nat-finite* **by** *blast*

```

have bij-betw (partitions-of  $A\ B$ ) ( $\{f \in A \rightarrow_E B. \text{inj-on } f\ A\}$  // range-permutation
 $A\ B$ )  $\{P. \text{partition-on } A\ P \wedge \text{card } P \leq \text{card } B \wedge (\forall X \in P. \text{card } X = 1)\}$ 
  using  $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$  by (rule bij-betw-partitions-of)
  from this have card ( $\{f \in A \rightarrow_E B. \text{inj-on } f\ A\}$  // range-permutation  $A\ B$ ) =
card  $\{P. \text{partition-on } A\ P \wedge \text{card } P \leq \text{card } B \wedge (\forall X \in P. \text{card } X = 1)\}$ 
  by (rule bij-betw-same-card)
  also have card  $\{P. \text{partition-on } A\ P \wedge \text{card } P \leq \text{card } B \wedge (\forall X \in P. \text{card } X =$ 
 $1)\}$  = iverson ( $\text{card } A \leq \text{card } B$ )
  using  $\langle \text{finite } A \rangle$  by (rule card-partition-on-size1-eq-iverson)
  finally show ?thesis .
qed

end

```

11 Surjections from A to B up to a Permutation on B

```

theory Twelvefold-Way-Entry9
imports Twelvefold-Way-Entry7
begin

```

11.1 Properties for Bijections

```

lemma surjective-on-implies-card-eq:
  assumes  $f \text{ ' } A = B$ 
  shows card  $((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}) = \text{card } B$ 
proof –
  from  $\langle f \text{ ' } A = B \rangle$  have  $\{\} \notin (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B$  by auto
  from  $\langle f \text{ ' } A = B \rangle$  have inj-on  $(\lambda b. \{x \in A. f\ x = b\})\ B$  by (fastforce intro: inj-onI)
  have card  $((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}) = \text{card } ((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B)$ 
  using  $\langle \{\} \notin (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B \rangle$  by simp
  also have  $\dots = \text{card } B$ 
  using  $\langle \text{inj-on } (\lambda b. \{x \in A. f\ x = b\})\ B \rangle$  by (rule card-image)
  finally show ?thesis .
qed

```

```

lemma card-eq-implies-surjective-on:
  assumes finite  $B\ f \in A \rightarrow_E B$ 
  assumes card-eq: card  $((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}) = \text{card } B$ 
  shows  $f \text{ ' } A = B$ 
proof
  from  $\langle f \in A \rightarrow_E B \rangle$  show  $f \text{ ' } A \subseteq B$  by auto
next
  show  $B \subseteq f \text{ ' } A$ 
proof
  fix  $x$ 

```

```

assume  $x \in B$ 
have  $\{\} \notin (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B$ 
proof (cases card  $B \geq 1$ )
  assume  $\neg \text{card } B \geq 1$ 
  from this have  $\text{card } B = 0$  by simp
  from this  $\langle \text{finite } B \rangle$  have  $B = \{\}$  by simp
  from this show ?thesis by simp
next
  assume  $\text{card } B \geq 1$ 
  show ?thesis
  proof (rule ccontr)
    assume  $\neg \{\} \notin (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B$ 
    from this have  $\{\} \in (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B$  by simp
    moreover have  $\text{card } ((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B) \leq \text{card } B$ 
      using  $\langle \text{finite } B \rangle$  card-image-le by blast
    moreover have  $\text{finite } ((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B)$ 
      using  $\langle \text{finite } B \rangle$  by auto
    ultimately have  $\text{card } ((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}) \leq \text{card } B - 1$ 
      by (auto simp add: card-Diff-singleton)
    from this card-eq  $\langle \text{card } B \geq 1 \rangle$  show False by auto
  qed
qed
from this  $\langle x \in B \rangle$  show  $x \in f \text{ ' } A$  by force
qed
qed

lemma card-partitions-of:
  assumes  $F \in (A \rightarrow_E B) // \text{range-permutation } A\ B$ 
  assumes  $\text{univ } (\lambda f. f \text{ ' } A = B)\ F$ 
  shows  $\text{card } (\text{partitions-of } A\ B\ F) = \text{card } B$ 
proof -
  from  $\langle F \in (A \rightarrow_E B) // \text{range-permutation } A\ B \rangle$  obtain  $f$  where  $f \in A \rightarrow_E B$ 
  and  $F\text{-eq: } F = \text{range-permutation } A\ B \text{ " } \{f\}$  using quotientE by blast
  from this  $\langle \text{univ } (\lambda f. f \text{ ' } A = B)\ F \rangle$  have  $f \text{ ' } A = B$ 
  using univ-commute'[OF equiv-range-permutation surj-on-respects-range-permutation
   $\langle f \in A \rightarrow_E B \rangle]$  by simp
  have  $\text{card } (\text{partitions-of } A\ B\ F) = \text{card } (\text{univ } (\lambda f. (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\})\ F)$ 
  unfolding partitions-of-def ..
  also have  $\dots = \text{card } (\text{univ } (\lambda f. (\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\}) (\text{range-permutation } A\ B \text{ " } \{f\}))$ 
  unfolding F-eq ..
  also have  $\dots = \text{card } ((\lambda b. \{x \in A. f\ x = b\}) \text{ ' } B - \{\{\}\})$ 
  using equiv-range-permutation domain-partitions-respects-range-permutation  $\langle f \in A \rightarrow_E B \rangle$ 
  by (subst univ-commute') auto
  also from  $\langle f \text{ ' } A = B \rangle$  have  $\dots = \text{card } B$ 
  using surjective-on-implys-card-eq by auto

```

finally show ?thesis .
qed

lemma functions-of-is-surj-on:
 assumes finite A finite B
 assumes partition-on A P card P = card B
 shows univ ($\lambda f. f \text{ ‘ } A = B$) (functions-of P A B)
 proof –
 have functions-of P A B $\in (A \rightarrow_E B)$ // range-permutation A B
 using functions-of ⟨finite A⟩ ⟨finite B⟩ ⟨partition-on A P⟩ ⟨card P = card B⟩
 by fastforce
 from this obtain f where eq-f: functions-of P A B = range-permutation A B
 “ {f} and $f \in A \rightarrow_E B$
 using quotientE by blast
 from eq-f have $f \in \text{functions-of } P \text{ A B}$
 using ⟨ $f \in A \rightarrow_E B$ ⟩ equiv-range-permutation equiv-class-self by fastforce
 from ⟨ $f \in \text{functions-of } P \text{ A B}$ ⟩ have eq: ($\lambda b. \{x \in A. f x = b\}$) ‘ $B - \{\{\}\} = P$
 unfolding functions-of-def by auto
 from this have card (($\lambda b. \{x \in A. f x = b\}$) ‘ $B - \{\{\}\}$) = card B
 using ⟨card P = card B⟩ by simp
 from ⟨finite B⟩ ⟨ $f \in A \rightarrow_E B$ ⟩ this have $f \text{ ‘ } A = B$
 using card-eq-implies-surjective-on by blast
 from this show ?thesis
 unfolding eq-f using equiv-range-permutation surj-on-respects-range-permutation
 ⟨ $f \in A \rightarrow_E B$ ⟩
 by (subst univ-commute’) assumption+
 qed

11.2 Bijections

lemma bij-betw-partitions-of:
 assumes finite A finite B
 shows bij-betw (partitions-of A B) ($\{f \in A \rightarrow_E B. f \text{ ‘ } A = B\}$ // range-permutation A B) {P. partition-on A P \wedge card P = card B}
 proof (rule bij-betw-byWitness[where f’= $\lambda P. \text{functions-of } P \text{ A B}$])
 have quotient-eq: $\{f \in A \rightarrow_E B. f \text{ ‘ } A = B\}$ // range-permutation A B = {F \in ((A \rightarrow_E B) // range-permutation A B). univ ($\lambda f. f \text{ ‘ } A = B$) F}
 using equiv-range-permutation[of A B] surj-on-respects-range-permutation[of A B] by (simp only: univ-preserves-predicate)
 show $\forall F \in \{f \in A \rightarrow_E B. f \text{ ‘ } A = B\}$ // range-permutation A B. functions-of (partitions-of A B F) A B = F
 using ⟨finite B⟩ by (simp add: functions-of-partitions-of quotient-eq)
 show $\forall P \in \{P. \text{partition-on } A \text{ P} \wedge \text{card } P = \text{card } B\}. \text{partitions-of } A \text{ B (functions-of } P \text{ A B)} = P$
 using ⟨finite A⟩ ⟨finite B⟩ by (auto simp add: partitions-of-functions-of)
 show partitions-of A B ‘ ($\{f \in A \rightarrow_E B. f \text{ ‘ } A = B\}$ // range-permutation A B)
 $\subseteq \{P. \text{partition-on } A \text{ P} \wedge \text{card } P = \text{card } B\}$
 using ⟨finite B⟩ quotient-eq card-partitions-of partitions-of by fastforce
 show ($\lambda P. \text{functions-of } P \text{ A B}$) ‘ {P. partition-on A P \wedge card P = card B} \subseteq


```

{f ∈ A →E B. f ‘ A = B} // range-permutation A B
  using ⟨finite A⟩ ⟨finite B⟩ by (auto simp add: quotient-eq intro: functions-of
functions-of-is-surj-on)
qed

```

11.3 Cardinality

lemma *card-surjective-functions-range-permutation:*

```

  assumes finite A finite B
  shows card ({f ∈ A →E B. f ‘ A = B} // range-permutation A B) = Stirling
(card A) (card B)
proof -
  have bij-betw (partitions-of A B) ({f ∈ A →E B. f ‘ A = B} // range-permutation
A B) {P. partition-on A P ∧ card P = card B}
  using ⟨finite A⟩ ⟨finite B⟩ by (rule bij-betw-partitions-of)
  from this have card ({f ∈ A →E B. f ‘ A = B} // range-permutation A B) =
card {P. partition-on A P ∧ card P = card B}
  by (rule bij-betw-same-card)
  also have card {P. partition-on A P ∧ card P = card B} = Stirling (card A)
(card B)
  using ⟨finite A⟩ by (rule card-partition-on)
  finally show ?thesis .
qed

```

end

12 Surjections from A to B

theory *Twelvefold-Way-Entry3*

imports

Twelvefold-Way-Entry9

begin

lemma *card-of-equiv-class:*

```

  assumes finite B
  assumes F ∈ {f ∈ A →E B. f ‘ A = B} // range-permutation A B
  shows card F = fact (card B)
proof -
  from ⟨F ∈ {f ∈ A →E B. f ‘ A = B} // range-permutation A B⟩ obtain f
where
  f ∈ A →E B and f ‘ A = B
  and F-eq: F = range-permutation A B “ {f} using quotientE by blast
  have set-eq: range-permutation A B “ {f} = (λp x. if x ∈ A then p (f x) else
undefined) ‘ {p. p permutes B}
  proof
    show range-permutation A B “ {f} ⊆ (λp x. if x ∈ A then p (f x) else undefined)
‘ {p. p permutes B}
    proof
      fix f'

```

```

assume  $f' \in \text{range-permutation } A \ B \text{ “}\{f\}$ 
from this obtain  $p$  where  $p$  permutes  $B \ \forall x \in A. f \ x = p \ (f' \ x)$ 
  unfolding range-permutation-def by auto
from  $\langle f' \in \text{range-permutation } A \ B \text{ “}\{f\} \rangle$  have  $f' \in A \rightarrow_E B$ 
  unfolding range-permutation-def by auto
have  $f' = (\lambda x. \text{if } x \in A \text{ then } \text{inv } p \ (f \ x) \text{ else undefined})$ 
proof
  fix  $x$ 
  show  $f' \ x = (\text{if } x \in A \text{ then } \text{inv } p \ (f \ x) \text{ else undefined})$ 
    using  $\langle f' \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle \langle \forall x \in A. f \ x = p \ (f' \ x) \rangle$ 
     $\langle p \text{ permutes } B \rangle \text{permutes-inverses}(2)$  by fastforce
qed
  moreover have  $\text{inv } p \text{ permutes } B$  using  $\langle p \text{ permutes } B \rangle$  by (simp add:
permutes-inv)
  ultimately show  $f' \in (\lambda p. (\lambda x. \text{if } x \in A \text{ then } p \ (f \ x) \text{ else undefined})) \text{ “}\{p. p \text{ permutes } B\}$ 
    by auto
qed
next
  show  $(\lambda p \ x. \text{if } x \in A \text{ then } p \ (f \ x) \text{ else undefined}) \text{ “}\{p. p \text{ permutes } B\} \subseteq$ 
range-permutation  $A \ B \text{ “}\{f\}$ 
proof
  fix  $f'$ 
  assume  $f' \in (\lambda p \ x. \text{if } x \in A \text{ then } p \ (f \ x) \text{ else undefined}) \text{ “}\{p. p \text{ permutes } B\}$ 
  from this obtain  $p$  where  $p$  permutes  $B$  and  $f'\text{-eq: } f' = (\lambda x. \text{if } x \in A \text{ then } p \ (f \ x) \text{ else undefined})$  by auto
  from this have  $f' \in A \rightarrow_E B$ 
    using  $\langle f' \in A \rightarrow_E B \rangle$  permutes-in-image by fastforce
    moreover have  $\text{inv } p \text{ permutes } B$  using  $\langle p \text{ permutes } B \rangle$  by (simp add:
permutes-inv)
    moreover have  $\forall x \in A. f \ x = \text{inv } p \ (f' \ x)$ 
      using  $\langle f' \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle f'\text{-eq}$ 
       $\langle p \text{ permutes } B \rangle \text{permutes-inverses}(2)$  by fastforce
    ultimately show  $f' \in \text{range-permutation } A \ B \text{ “}\{f\}$ 
      using  $\langle f' \in A \rightarrow_E B \rangle$  unfolding range-permutation-def by auto
qed
qed
have inj-on  $(\lambda p \ x. \text{if } x \in A \text{ then } p \ (f \ x) \text{ else undefined}) \ \{p. p \text{ permutes } B\}$ 
proof (rule inj-onI)
  fix  $p \ p'$ 
  assume  $p \in \{p. p \text{ permutes } B\} \ p' \in \{p. p \text{ permutes } B\}$ 
  and eq:  $(\lambda x. \text{if } x \in A \text{ then } p \ (f \ x) \text{ else undefined}) = (\lambda x. \text{if } x \in A \text{ then } p' \ (f \ x) \text{ else undefined})$ 
  {
    fix  $x$ 
    have  $p \ x = p' \ x$ 
    proof cases
      assume  $x \in B$ 
      from this obtain  $y$  where  $y \in A$  and  $x = f \ y$ 

```

```

      using ⟨f ‘ A = B⟩ by blast
    from eq this have p (f y) = p' (f y) by meson
    from this ⟨x = f y⟩ show p x = p' x by simp
  next
    assume x ∉ B
    from this show p x = p' x
      using ⟨p ∈ {p. p permutes B}⟩ ⟨p' ∈ {p. p permutes B}⟩
      by (simp add: permutes-def)
    qed
  }
  from this show p = p' by auto
  qed
  have card F = card ((λp x. if x ∈ A then p (f x) else undefined) ‘ {p. p permutes B})
  unfolding F-eq set-eq ..
  also have ... = card {p. p permutes B}
    using ⟨inj-on (λp x. if x ∈ A then p (f x) else undefined) {p. p permutes B}⟩
    by (simp add: card-image)
  also have ... = fact (card B)
    using ⟨finite B⟩ by (simp add: card-permutations)
  finally show ?thesis .
  qed

lemma card-extensional-funcset-surj-on:
  assumes finite A finite B
  shows card {f ∈ A →E B. f ‘ A = B} = fact (card B) * Stirling (card A) (card B) (is card ?F = -)
  proof -
    have card ?F = fact (card B) * card (?F // range-permutation A B)
      using ⟨finite B⟩
      by (simp only: card-equiv-class-restricted-same-size[OF equiv-range-permutation surj-on-respects-range-permutation card-of-equiv-class])
    also have ... = fact (card B) * Stirling (card A) (card B)
      using ⟨finite A⟩ ⟨finite B⟩
      by (simp only: card-surjective-functions-range-permutation)
    finally show ?thesis .
  qed

end

```

13 Functions from A to B up to a Permutation on A and B

```

theory Twelvefold-Way-Entry10
imports Equiv-Relations-on-Functions
begin

```

13.1 Definition of Bijections

definition *number-partition-of* :: 'a set \Rightarrow 'b set \Rightarrow ('a \Rightarrow 'b) set \Rightarrow nat multiset where

number-partition-of A B F = univ (λf . image-mset (λX . card X) (mset-set ((λb . { $x \in A$. f x = b}) ' B - { $\{\{\}\}$ })) F

definition *functions-of* :: 'a set \Rightarrow 'b set \Rightarrow nat multiset \Rightarrow ('a \Rightarrow 'b) set where

functions-of A B N = { $f \in A \rightarrow_E B$. image-mset (λX . card X) (mset-set ((λb . { $x \in A$. f x = b}) ' B - { $\{\{\}\}$ })) = N}

13.2 Properties for Bijections

lemma *card-setsum-partition*:

assumes *finite* A *finite* B $f \in A \rightarrow_E B$

shows *sum card* ((λb . { $x \in A$. f x = b}) ' B - { $\{\{\}\}$ }) = *card* A

proof –

have *finite* ((λb . { $x \in A$. f x = b}) ' B - { $\{\{\}\}$ })

using <finite B> **by** blast

moreover have $\forall X \in (\lambda b$. { $x \in A$. f x = b}) ' B - { $\{\{\}\}$. *finite* X

using <finite A> **by** auto

moreover have $\bigcup ((\lambda b$. { $x \in A$. f x = b}) ' B - { $\{\{\}\}$ }) = A

using < $f \in A \rightarrow_E B$ > **by** auto

ultimately show ?thesis

by (subst card-Union-disjoint[symmetric]) (auto simp: pairwise-def disjnt-def)

qed

lemma *number-partition-of*:

assumes *finite* A *finite* B

assumes $F \in (A \rightarrow_E B)$ // *domain-and-range-permutation* A B

shows *number-partition* (card A) (*number-partition-of* A B F)

and *size* (*number-partition-of* A B F) \leq *card* B

proof –

from < $F \in (A \rightarrow_E B)$ // *domain-and-range-permutation* A B> **obtain** f **where** $f \in A \rightarrow_E B$

and *F-eq*: $F = \text{domain-and-range-permutation } A \ B \ \text{"}\{f\}$ **using** quotientE **by** blast

have *number-partition-of-eq*: *number-partition-of* A B F = *image-mset* card (mset-set ((λb . { $x \in A$. f x = b}) ' B - { $\{\{\}\}$ }))

proof –

have *number-partition-of* A B F = univ (λf . *image-mset* card (mset-set ((λb . { $x \in A$. f x = b}) ' B - { $\{\{\}\}$ })) F

unfolding *number-partition-of-def* ..

also have ... = univ (λf . *image-mset* card (mset-set ((λb . { $x \in A$. f x = b}) ' B - { $\{\{\}\}$ })) (domain-and-range-permutation A B " {f})

unfolding *F-eq* ..

also have ... = *image-mset* card (mset-set ((λb . { $x \in A$. f x = b}) ' B - { $\{\{\}\}$ }))

using <finite B> *equiv-domain-and-range-permutation multiset-of-partition-cards-respects-domain-and-range*

```

  ⟨f ∈ A →E B⟩
    by (subst univ-commute') auto
    finally show ?thesis .
qed
show number-partition (card A) (number-partition-of A B F)
proof -
  have sum-mset (number-partition-of A B F) = card A
    using number-partition-of-eq ⟨finite A⟩ ⟨finite B⟩ ⟨f ∈ A →E B⟩
    by (simp only: sum-unfold-sum-mset[symmetric] card-setsum-partition)
  moreover have 0 ∉ # number-partition-of A B F
  proof -
    have ∀ X ∈ (λb. {x ∈ A. f x = b}) ‘ B. finite X
      using ⟨finite A⟩ by simp
    from this have ∀ X ∈ (λb. {x ∈ A. f x = b}) ‘ B - {{}}. card X ≠ 0 by
auto
    from this show ?thesis
      using number-partition-of-eq ⟨finite B⟩ by (simp add: image-iff)
  qed
  ultimately show ?thesis unfolding number-partition-def by simp
qed
show size (number-partition-of A B F) ≤ card B
  using number-partition-of-eq ⟨finite A⟩ ⟨finite B⟩
  by (metis (no-types, lifting) card-Diff1-le card-image-le finite-imageI le-trans
size-image-mset size-mset-set)
qed

lemma functions-of:
  assumes finite A finite B
  assumes number-partition (card A) N
  assumes size N ≤ card B
  shows functions-of A B N ∈ (A →E B) // domain-and-range-permutation A B
proof -
  obtain f where f ∈ A →E B and eq-N: image-mset (λX. card X) (mset-set
(((λb. {x ∈ A. f x = b}))) ‘ B - {{}})) = N
    using obtain-extensional-function-from-number-partition ⟨finite A⟩ ⟨finite B⟩
  ⟨number-partition (card A) N⟩ ⟨size N ≤ card B⟩ by blast
  have functions-of A B N = (domain-and-range-permutation A B) “ {f}
  proof
    show functions-of A B N ⊆ domain-and-range-permutation A B “ {f}
  proof
    fix f'
    assume f' ∈ functions-of A B N
    from this have eq-N': N = image-mset (λX. card X) (mset-set (((λb. {x ∈
A. f' x = b}))) ‘ B - {{}}))
      and f' ∈ A →E B
    unfolding functions-of-def by auto
    from ⟨finite A⟩ ⟨finite B⟩ ⟨f ∈ A →E B⟩ ⟨f' ∈ A →E B⟩
    obtain pA pB where pA permutes A pB permutes B ∀ x ∈ A. f x = pB (f' (pA
x))

```

```

    using eq-N eq-N' multiset-of-partition-cards-eq-implies-permutes[of A B f f']
  by blast
  from this show f' ∈ domain-and-range-permutation A B “ {f}
    using ⟨f ∈ A →E B⟩ ⟨f' ∈ A →E B⟩
    unfolding domain-and-range-permutation-def by auto
  qed
next
show domain-and-range-permutation A B “ {f} ⊆ functions-of A B N
proof
  fix f'
  assume f' ∈ domain-and-range-permutation A B “ {f}
  from this have in-equiv-relation: (f, f') ∈ domain-and-range-permutation A
B by auto
  from eq-N ⟨finite B⟩ have image-mset (λX. card X) (mset-set ((λb. {x ∈
A. f' x = b})) ‘ B - {f})) = N
  using congruentD[OF multiset-of-partition-cards-respects-domain-and-range-permutation
in-equiv-relation]
  by metis
  moreover from ⟨(f, f') ∈ domain-and-range-permutation A B⟩ have f' ∈ A
→E B
  unfolding domain-and-range-permutation-def by auto
  ultimately show f' ∈ functions-of A B N
  unfolding functions-of-def by auto
qed
qed
from this ⟨f ∈ A →E B⟩ show ?thesis by (auto intro: quotientI)
qed

```

lemma functions-of-number-partition-of:

```

  assumes finite A finite B
  assumes F ∈ (A →E B) // domain-and-range-permutation A B
  shows functions-of A B (number-partition-of A B F) = F
proof -
  from ⟨F ∈ (A →E B) // domain-and-range-permutation A B⟩ obtain f where
f ∈ A →E B
  and F-eq: F = domain-and-range-permutation A B “ {f} using quotientE by
blast
  have number-partition-of A B F = univ (λf. image-mset card (mset-set ((λb. {x
∈ A. f x = b}) ‘ B - {f}))) F
  unfolding number-partition-of-def ..
  also have ... = univ (λf. image-mset card (mset-set ((λb. {x ∈ A. f x = b}) ‘
B - {f}))) (domain-and-range-permutation A B “ {f})
  unfolding F-eq ..
  also have ... = image-mset card (mset-set ((λb. {x ∈ A. f x = b}) ‘ B - {f}))
  using ⟨finite B⟩
  using equiv-domain-and-range-permutation multiset-of-partition-cards-respects-domain-and-range-permutati
⟨f ∈ A →E B⟩
  by (subst univ-commute') auto
  finally have number-partition-of-eq: number-partition-of A B F = image-mset

```

```

card (mset-set ((λb. {x ∈ A. f x = b}) ‘ B - {{}))) .
show ?thesis
proof
  show functions-of A B (number-partition-of A B F) ⊆ F
  proof
    fix f'
    assume f' ∈ functions-of A B (number-partition-of A B F)
    from this have f' ∈ A →E B
    and eq: image-mset card (mset-set ((λb. {x ∈ A. f' x = b}) ‘ B - {{})))
= image-mset card (mset-set ((λb. {x ∈ A. f x = b}) ‘ B - {{})))
    unfolding functions-of-def by (auto simp add: number-partition-of-eq)
    note ⟨f ∈ A →E B⟩ ⟨f' ∈ A →E B⟩
    moreover obtain pA pB where pA permutes A pB permutes B ∀ x ∈ A. f x
= pB (f' (pA x))
    using ⟨finite A⟩ ⟨finite B⟩ ⟨f ∈ A →E B⟩ ⟨f' ∈ A →E B⟩ eq
    multiset-of-partition-cards-eq-implies-permutes[of A B f f']
    by metis
    ultimately show f' ∈ F
    unfolding F-eq domain-and-range-permutation-def by auto
  qed
next
show F ⊆ functions-of A B (number-partition-of A B F)
proof
  fix f'
  assume f' ∈ F
  from ⟨f' ∈ F⟩ obtain pA pB where pA permutes A pB permutes B ∀ x ∈ A.
f x = pB (f' (pA x))
    unfolding F-eq domain-and-range-permutation-def by auto
    have eq: image-mset card (mset-set ((λb. {x ∈ A. f x = b}) ‘ B - {{}))) =
image-mset card (mset-set ((λb. {x ∈ A. f' x = b}) ‘ B - {{})))
    proof -
      have (λb. {x ∈ A. f x = b}) ‘ B = (λb. {x ∈ A. pB (f' (pA x)) = b}) ‘ B
      using ⟨∀ x ∈ A. f x = pB (f' (pA x))⟩ by auto
      from this have image-mset card (mset-set ((λb. {x ∈ A. f x = b}) ‘ B -
{{}))) =
      image-mset card (mset-set ((λb. {x ∈ A. pB (f' (pA x)) = b}) ‘ B - {{})))
    by simp
    also have ... = image-mset card (mset-set ((λb. {x ∈ A. f' x = b}) ‘ B -
{{})))
    using ⟨pA permutes A⟩ ⟨pB permutes B⟩ permutes-implies-multiset-of-partition-cards-eq
  by blast
  finally show ?thesis .
qed
moreover from ⟨f' ∈ F⟩ have f' ∈ A →E B
  unfolding F-eq domain-and-range-permutation-def by auto
  ultimately show f' ∈ functions-of A B (number-partition-of A B F)
  unfolding functions-of-def number-partition-of-eq by auto
qed
qed

```

qed

lemma *number-partition-of-functions-of*:

assumes *finite A finite B*

assumes *number-partition (card A) N size N ≤ card B*

shows *number-partition-of A B (functions-of A B N) = N*

proof –

from *assms* **have** *functions-of A B N ∈ (A →_E B) // domain-and-range-permutation A B*

using *functions-of assms by fastforce*

from *this* **obtain** *f where f ∈ A →_E B and functions-of A B N = domain-and-range-permutation A B “{f}”*

by *(meson quotientE)*

from *this* **have** *f ∈ functions-of A B N*

using *equiv-domain-and-range-permutation equiv-class-self by fastforce*

have *number-partition-of A B (functions-of A B N) = univ (λf. image-mset card (mset-set ((λb. {x ∈ A. f x = b}) ‘B - {f}’))) (functions-of A B N)*

unfolding *number-partition-of-def ..*

also have *... = univ (λf. image-mset card (mset-set ((λb. {x ∈ A. f x = b}) ‘B - {f}’))) (domain-and-range-permutation A B “{f}”)*

unfolding *‘functions-of A B N = domain-and-range-permutation A B “{f}”*

..

also have *... = image-mset card (mset-set ((λb. {x ∈ A. f x = b}) ‘B - {f}’))*

using *‘finite B’ ‘f ∈ A →_E B’ equiv-domain-and-range-permutation*

multiset-of-partition-cards-respects-domain-and-range-permutation

by *(subst univ-commute’) auto*

also have *image-mset card (mset-set ((λb. {x ∈ A. f x = b}) ‘B - {f}’)) = N*

using *‘f ∈ functions-of A B N’ unfolding functions-of-def by simp*

finally show *?thesis .*

qed

13.3 Bijections

lemma *bij-betw-number-partition-of*:

assumes *finite A finite B*

shows *bij-betw (number-partition-of A B) ((A →_E B) // domain-and-range-permutation A B) {N. number-partition (card A) N ∧ size N ≤ card B}*

proof *(rule bij-betw-byWitness[where f’=λM. functions-of A B M])*

show *∀ F ∈ (A →_E B) // domain-and-range-permutation A B. functions-of A B (number-partition-of A B F) = F*

using *‘finite A’ ‘finite B’ by (auto simp add: functions-of-number-partition-of)*

show *∀ N ∈ {N. number-partition (card A) N ∧ size N ≤ card B}. number-partition-of A B (functions-of A B N) = N*

using *‘finite A’ ‘finite B’ by (auto simp add: number-partition-of-functions-of)*

show *number-partition-of A B ‘((A →_E B) // domain-and-range-permutation A B) ⊆ {N. number-partition (card A) N ∧ size N ≤ card B}*

using *number-partition-of[of A B] ‘finite A’ ‘finite B’ by auto*

show *functions-of A B ‘{N. number-partition (card A) N ∧ size N ≤ card B} ⊆ (A →_E B) // domain-and-range-permutation A B*

using *functions-of* $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ by *blast*
qed

13.4 Cardinality

lemma *card-domain-and-range-permutation*:

assumes *finite A finite B*
shows $\text{card } ((A \rightarrow_E B) // \text{domain-and-range-permutation } A \ B) = \text{Partition}$
 $(\text{card } A + \text{card } B) (\text{card } B)$
proof –
have *bij-betw (number-partition-of A B) ((A →_E B) // domain-and-range-permutation*
A B) {N. number-partition (card A) N ∧ size N ≤ card B}
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ by (rule *bij-betw-number-partition-of*)
from this have $\text{card } ((A \rightarrow_E B) // \text{domain-and-range-permutation } A \ B) = \text{card}$
 $\{N. \text{number-partition } (\text{card } A) \ N \wedge \text{size } N \leq \text{card } B\}$
by (rule *bij-betw-same-card*)
also have $\text{card } \{N. \text{number-partition } (\text{card } A) \ N \wedge \text{size } N \leq \text{card } B\} = \text{Partition}$
 $(\text{card } A + \text{card } B) (\text{card } B)$
by (rule *card-number-partitions-with-atmost-k-parts*)
finally show ?thesis .
qed

end

14 Injections from A to B up to a permutation on A and B

theory *Twelvefold-Way-Entry11*
imports *Twelvefold-Way-Entry10*
begin

14.1 Properties for Bijections

lemma *all-one-implies-inj-on*:

assumes *finite A finite B*
assumes $\forall n. n \in \# \ N \longrightarrow n = 1 \ \text{number-partition } (\text{card } A) \ N \ \text{size } N \leq \text{card } B$
assumes $f \in \text{functions-of } A \ B \ N$
shows *inj-on f A*
proof –
from $\langle f \in \text{functions-of } A \ B \ N \rangle$ have $f \in A \rightarrow_E B$
and $N = \text{image-mset card (mset-set } ((\lambda b. \{x \in A. f \ x = b\}) \text{ ‘ } B - \{\{\}\}))$
unfolding *functions-of-def* by *auto*
from this $\langle \forall n. n \in \# \ N \longrightarrow n = 1 \rangle$ have *parts*: $\forall b \in B. \text{card } \{x \in A. f \ x = b\}$
 $= 1 \vee \{x \in A. f \ x = b\} = \{\}$
using $\langle \text{finite } B \rangle$ by *auto*
show *inj-on f A*
proof
fix $x \ y$
assume $a: x \in A \ y \in A \ f \ x = f \ y$

```

from  $\langle f \in A \rightarrow_E B \rangle \langle x \in A \rangle$  have  $f x \in B$  by auto
from  $a$  have  $1: x \in \{x' \in A. f x' = f x\} \ y \in \{x' \in A. f x' = f x\}$  by auto
from this have  $2: \text{card } \{x' \in A. f x' = f x\} = 1$ 
  using parts  $\langle f x \in B \rangle$  by blast
from this have is-singleton  $\{x' \in A. f x' = f x\}$ 
  by (simp add: is-singleton-altdef)
from  $1$  this show  $x = y$ 
  by (metis is-singletonE singletonD)
qed
qed

```

lemma *inj-on-implies-all-one*:

```

assumes finite A finite B
assumes  $F \in (A \rightarrow_E B)$  // domain-and-range-permutation A B
assumes univ  $(\lambda f. \text{inj-on } f A) F$ 
shows  $\forall n. n \in \# \text{ number-partition-of } A B F \longrightarrow n = 1$ 
proof -
  from  $\langle F \in (A \rightarrow_E B) \text{ // domain-and-range-permutation } A B \rangle$  obtain  $f$  where
     $f \in A \rightarrow_E B$ 
  and  $F\text{-eq}: F = \text{domain-and-range-permutation } A B \text{ ``}\{f\}$  using quotientE by
    blast
  have number-partition-of  $A B F = \text{univ } (\lambda f. \text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f x = b\}) \text{ `` } B - \{\{\}\}))) F$ 
  unfolding number-partition-of-def ..
  also have  $\dots = \text{univ } (\lambda f. \text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f x = b\}) \text{ `` } B - \{\{\}\}))) (\text{domain-and-range-permutation } A B \text{ `` } \{f\})$ 
  unfolding  $F\text{-eq}$  ..
  also have  $\dots = \text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f x = b\}) \text{ `` } B - \{\{\}\})))$ 
  using  $\langle \text{finite } B \rangle$  equiv-domain-and-range-permutation multiset-of-partition-cards-respects-domain-and-range-
     $\langle f \in A \rightarrow_E B \rangle$ 
  by (subst univ-commute') auto
  finally have  $\text{eq: number-partition-of } A B F = \text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f x = b\}) \text{ `` } B - \{\{\}\})))$  .
  from iffD1 [OF univ-commute', OF equiv-domain-and-range-permutation, OF inj-on-respects-domain-and-range-permutation, OF  $\langle f \in A \rightarrow_E B \rangle$ ]
    assms(4) have inj-on  $f A$  by (simp add: F-eq)
  have  $\forall n. n \in \# \text{ image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f x = b\}) \text{ `` } B - \{\{\}\}))) \longrightarrow n = 1$ 
proof -
  have  $\forall b \in B. \text{card } \{x \in A. f x = b\} = 1 \vee \{x \in A. f x = b\} = \{\}$ 
proof
  fix  $b$ 
  assume  $b \in B$ 
  show  $\text{card } \{x \in A. f x = b\} = 1 \vee \{x \in A. f x = b\} = \{\}$ 
proof (cases  $b \in f \text{ `` } A$ )
    assume  $b \in f \text{ `` } A$ 
    from  $\langle \text{inj-on } f A \rangle$  this have is-singleton  $\{x \in A. f x = b\}$ 
    by (auto simp add: inj-on-eq-iff intro: is-singletonI')
    from this have  $\text{card } \{x \in A. f x = b\} = 1$ 

```

```

      by (subst is-singleton-altdef[symmetric])
    from this show ?thesis ..
  next
    assume  $b \notin f \text{ `` } A$ 
    from this have  $\{x \in A. f\ x = b\} = \{\}$  by auto
    from this show ?thesis ..
  qed
qed
from this show ?thesis
  using ⟨finite B⟩ by auto
qed
from this show  $\forall n. n \in \# \text{ number-partition-of } A\ B\ F \longrightarrow n = 1$ 
  unfolding eq by auto
qed

```

lemma *functions-of-is-inj-on:*

```

  assumes finite A finite B
  assumes  $\forall n. n \in \# N \longrightarrow n = 1 \text{ number-partition } (\text{card } A)\ N \text{ size } N \leq \text{card } B$ 
  shows univ  $(\lambda f. \text{inj-on } f\ A)$   $(\text{functions-of } A\ B\ N)$ 
proof -
  have functions-of  $A\ B\ N \in (A \rightarrow_E B)$  // domain-and-range-permutation A B
    using assms functions-of by auto
  from this obtain f where eq-f: functions-of  $A\ B\ N = \text{domain-and-range-permutation}$ 
     $A\ B\ \text{ `` } \{f\}$  and  $f \in A \rightarrow_E B$ 
    using quotientE by blast
  from eq-f have  $f \in \text{functions-of } A\ B\ N$ 
    using ⟨ $f \in A \rightarrow_E B$ ⟩ equiv-domain-and-range-permutation equiv-class-self by
  fastforce
  have inj-on f A
    using ⟨ $f \in \text{functions-of } A\ B\ N$ ⟩ assms all-one-implies-inj-on by blast
  from this show ?thesis
    unfolding eq-f using equiv-domain-and-range-permutation inj-on-respects-domain-and-range-permutation
    ⟨ $f \in A \rightarrow_E B$ ⟩
    by (subst univ-commute') assumption+
qed

```

14.2 Bijections

lemma *bij-betw-number-partition-of:*

```

  assumes finite A finite B
  shows bij-betw (number-partition-of A B)  $(\{f \in A \rightarrow_E B. \text{inj-on } f\ A\}$  // domain-and-range-permutation A B)
 $\{N. (\forall n. n \in \# N \longrightarrow n = 1) \wedge \text{number-partition } (\text{card } A)\ N \wedge \text{size } N \leq \text{card } B\}$ 
proof (rule bij-betw-byWitness[where f'=functions-of A B])
  have quotient-eq:  $\{f \in A \rightarrow_E B. \text{inj-on } f\ A\}$  // domain-and-range-permutation
     $A\ B = \{F \in ((A \rightarrow_E B) // \text{domain-and-range-permutation } A\ B). \text{univ } (\lambda f. \text{inj-on } f\ A)\ F\}$ 
    using equiv-domain-and-range-permutation[of A B] inj-on-respects-domain-and-range-permutation[of
    A B] by (simp only: univ-preserves-predicate)

```

show $\forall F \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-and-range-permutation } A \ B.$
 $\text{functions-of } A \ B \ (\text{number-partition-of } A \ B \ F) = F$
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ **by** (auto simp only: quotient-eq functions-of-number-partition-of)
show $\forall N \in \{N. (\forall n. n \in \# \ N \longrightarrow n = 1) \wedge \text{number-partition } (\text{card } A) \ N \wedge \text{size}$
 $N \leq \text{card } B\}. \text{number-partition-of } A \ B \ (\text{functions-of } A \ B \ N) = N$
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ $\text{number-partition-of-functions-of}$ **by** auto
show $\text{number-partition-of } A \ B \ ' (\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-and-range-permutation}$
 $A \ B)$
 $\subseteq \{N. (\forall n. n \in \# \ N \longrightarrow n = 1) \wedge \text{number-partition } (\text{card } A) \ N \wedge \text{size } N \leq$
 $\text{card } B\}$
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$
by (auto simp add: quotient-eq number-partition-of inj-on-implies-all-one simp
del: One-nat-def)
show $\text{functions-of } A \ B \ ' \{N. (\forall n. n \in \# \ N \longrightarrow n = 1) \wedge \text{number-partition } (\text{card}$
 $A) \ N \wedge \text{size } N \leq \text{card } B\}$
 $\subseteq \{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-and-range-permutation } A \ B$
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ **by** (auto simp add: quotient-eq intro: functions-of
functions-of-is-inj-on)
qed

lemma *bij-betw-functions-of:*

assumes $\text{finite } A \ \text{finite } B$
shows $\text{bij-betw } (\text{functions-of } A \ B) \ \{N. (\forall n. n \in \# \ N \longrightarrow n = 1) \wedge \text{num-}$
 $\text{ber-partition } (\text{card } A) \ N \wedge \text{size } N \leq \text{card } B\} \ (\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} //$
 $\text{domain-and-range-permutation } A \ B)$
proof (rule *bij-betw-byWitness*[**where** $f' = \text{number-partition-of } A \ B$])
have $\text{quotient-eq: } \{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-and-range-permutation}$
 $A \ B = \{F \in ((A \rightarrow_E B) // \text{domain-and-range-permutation } A \ B). \text{univ } (\lambda f. \text{inj-on}$
 $f \ A) \ F\}$
using $\text{equiv-domain-and-range-permutation}[of \ A \ B] \ \text{inj-on-respects-domain-and-range-permutation}[of$
 $A \ B]$ **by** (simp only: univ-preserves-predicate)
show $\forall F \in \{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-and-range-permutation } A \ B.$
 $\text{functions-of } A \ B \ (\text{number-partition-of } A \ B \ F) = F$
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ **by** (auto simp only: quotient-eq functions-of-number-partition-of)
show $\forall N \in \{N. (\forall n. n \in \# \ N \longrightarrow n = 1) \wedge \text{number-partition } (\text{card } A) \ N \wedge \text{size}$
 $N \leq \text{card } B\}. \text{number-partition-of } A \ B \ (\text{functions-of } A \ B \ N) = N$
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ $\text{number-partition-of-functions-of}$ **by** auto
show $\text{number-partition-of } A \ B \ ' (\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-and-range-permutation}$
 $A \ B)$
 $\subseteq \{N. (\forall n. n \in \# \ N \longrightarrow n = 1) \wedge \text{number-partition } (\text{card } A) \ N \wedge \text{size } N \leq$
 $\text{card } B\}$
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$
by (auto simp add: quotient-eq number-partition-of inj-on-implies-all-one simp
del: One-nat-def)
show $\text{functions-of } A \ B \ ' \{N. (\forall n. n \in \# \ N \longrightarrow n = 1) \wedge \text{number-partition } (\text{card}$
 $A) \ N \wedge \text{size } N \leq \text{card } B\}$
 $\subseteq \{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-and-range-permutation } A \ B$
using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ **by** (auto simp add: quotient-eq intro: functions-of
functions-of-is-inj-on)

qed

14.3 Cardinality

lemma *card-injective-functions-domain-and-range-permutation:*

assumes *finite A finite B*

shows $\text{card } (\{f \in A \rightarrow_E B. \text{inj-on } f A\} // \text{domain-and-range-permutation } A B)$
 $= \text{iverson } (\text{card } A \leq \text{card } B)$

proof –

have *bij-betw (number-partition-of A B) ($\{f \in A \rightarrow_E B. \text{inj-on } f A\} // \text{domain-and-range-permutation } A B)$ $\{N. (\forall n. n \in \# N \longrightarrow n = 1) \wedge \text{number-partition (card A) } N \wedge \text{size } N \leq \text{card } B\}$*

using *$\langle \text{finite } A \rangle \langle \text{finite } B \rangle$ by (rule bij-betw-number-partition-of)*

from this have $\text{card } (\{f \in A \rightarrow_E B. \text{inj-on } f A\} // \text{domain-and-range-permutation } A B) = \text{card } \{N. (\forall n. n \in \# N \longrightarrow n = 1) \wedge \text{number-partition (card A) } N \wedge \text{size } N \leq \text{card } B\}$

by *(rule bij-betw-same-card)*

also have $\text{card } \{N. (\forall n. n \in \# N \longrightarrow n = 1) \wedge \text{number-partition (card A) } N \wedge \text{size } N \leq \text{card } B\} = \text{iverson } (\text{card } A \leq \text{card } B)$

by *(rule card-number-partitions-with-only-parts-1)*

finally show *?thesis .*

qed

end

15 Surjections from A to B up to a Permutation on A and B

theory *Twelvefold-Way-Entry12*

imports *Twelvefold-Way-Entry9 Twelvefold-Way-Entry10*

begin

15.1 Properties for Bijections

lemma *size-eq-card-implies-surj-on:*

assumes *finite A finite B*

assumes *size N = card B*

assumes *f ∈ functions-of A B N*

shows *f ‘ A = B*

proof –

from *$\langle f \in \text{functions-of } A B N \rangle$ have $f \in A \rightarrow_E B$ and*

$N = \text{image-mset card (mset-set } ((\lambda b. \{x \in A. f x = b\}) \text{ ‘ } B - \{\{\}\}))$

unfolding functions-of-def by auto

from this *size N = card B* **have** $\text{card } ((\lambda b. \{x \in A. f x = b\}) \text{ ‘ } B - \{\{\}\}) = \text{card } B$ **by** *simp*

from this *$\langle \text{finite } B \rangle \langle f \in A \rightarrow_E B \rangle$ show $f \text{ ‘ } A = B$*

using *card-eq-implies-surjective-on by blast*

qed

lemma *surj-on-implies-size-eq-card*:

assumes *finite A finite B*

assumes $F \in (A \rightarrow_E B)$ // *domain-and-range-permutation A B*

assumes *univ* $(\lambda f. f \text{ ‘ } A = B)$ *F*

shows *size* $(\text{number-partition-of } A \ B \ F) = \text{card } B$

proof –

from $\langle F \in (A \rightarrow_E B) \text{ // domain-and-range-permutation } A \ B \rangle$ **obtain** *f* **where**
 $f \in A \rightarrow_E B$

and *F-eq*: $F = \text{domain-and-range-permutation } A \ B \text{ ‘ ‘ } \{f\}$ **using** *quotientE* **by**
blast

have *number-partition-of* $A \ B \ F = \text{univ } (\lambda f. \text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f \ x = b\}) \text{ ‘ } B - \{\{\}\}))) \ F$

unfolding *number-partition-of-def* ..

also have $\dots = \text{univ } (\lambda f. \text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f \ x = b\}) \text{ ‘ } B - \{\{\}\}))) \ (\text{domain-and-range-permutation } A \ B \text{ ‘ ‘ } \{f\})$

unfolding *F-eq* ..

also have $\dots = \text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f \ x = b\}) \text{ ‘ } B - \{\{\}\})))$

using $\langle \text{finite } B \rangle$ *equiv-domain-and-range-permutation multiset-of-partition-cards-respects-domain-and-range-*
 $\langle f \in A \rightarrow_E B \rangle$

by *(subst univ-commute’)* *auto*

finally have *eq*: *number-partition-of* $A \ B \ F = \text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f \ x = b\}) \text{ ‘ } B - \{\{\}\})))$.

from *iffD1* [*OF univ-commute’*, *OF equiv-domain-and-range-permutation*, *OF*
surjective-respects-domain-and-range-permutation, *OF* $\langle f \in A \rightarrow_E B \rangle$]

assms (4) **have** $f \text{ ‘ } A = B$ **by** *(simp add: F-eq)*

have *size* $(\text{number-partition-of } A \ B \ F) = \text{size } (\text{image-mset card } (\text{mset-set } ((\lambda b. \{x \in A. f \ x = b\}) \text{ ‘ } B - \{\{\}\})))$

unfolding *eq* ..

also have $\dots = \text{card } ((\lambda b. \{x \in A. f \ x = b\}) \text{ ‘ } B - \{\{\}\})$ **by** *simp*

also from $\langle f \text{ ‘ } A = B \rangle$ **have** $\dots = \text{card } B$

using *surjective-on-implies-card-eq* **by** *auto*

finally show *?thesis* .

qed

lemma *functions-of-is-surj-on*:

assumes *finite A finite B*

assumes *number-partition* $(\text{card } A) \ N \ \text{size } N = \text{card } B$

shows *univ* $(\lambda f. f \text{ ‘ } A = B)$ $(\text{functions-of } A \ B \ N)$

proof –

have *functions-of* $A \ B \ N \in (A \rightarrow_E B)$ // *domain-and-range-permutation A B*

using *functions-of* $\langle \text{finite } A \rangle \langle \text{finite } B \rangle \langle \text{number-partition } (\text{card } A) \ N \rangle \langle \text{size } N = \text{card } B \rangle$

by *fastforce*

from this obtain *f* **where** *eq-f*: *functions-of* $A \ B \ N = \text{domain-and-range-permutation}$
 $A \ B \text{ ‘ ‘ } \{f\}$ **and** $f \in A \rightarrow_E B$

using *quotientE* **by** *blast*

from *eq-f* **have** $f \in \text{functions-of } A \ B \ N$

using $\langle f \in A \rightarrow_E B \rangle$ *equiv-domain-and-range-permutation equiv-class-self* **by**
fastforce

```

have f ' A = B
  using ⟨f ∈ functions-of A B N⟩ asms size-eq-card-implies-surj-on by blast
from this show ?thesis
  unfolding eq-f using equiv-domain-and-range-permutation surjective-respects-domain-and-range-permutation
  ⟨f ∈ A →E B⟩
  by (subst univ-commute') assumption+
qed

```

15.2 Bijections

lemma *bij-betw-number-partition-of*:

```

  assumes finite A finite B
  shows bij-betw (number-partition-of A B) ({f ∈ A →E B. f ' A = B} // domain-and-range-permutation A B) {N. number-partition (card A) N ∧ size N = card B}
proof (rule bij-betw-byWitness[where f'=functions-of A B])
  have quotient-eq: {f ∈ A →E B. f ' A = B} // domain-and-range-permutation A B = {F ∈ ((A →E B) // domain-and-range-permutation A B). univ (λf. f ' A = B) F}
  using equiv-domain-and-range-permutation[of A B] surjective-respects-domain-and-range-permutation[of A B] by (simp only: univ-preserves-predicate)
  show ∀ F ∈ {f ∈ A →E B. f ' A = B} // domain-and-range-permutation A B.
    functions-of A B (number-partition-of A B F) = F
  using ⟨finite A⟩ ⟨finite B⟩ by (auto simp only: quotient-eq functions-of-number-partition-of)
  show ∀ N ∈ {N. number-partition (card A) N ∧ size N = card B}. number-partition-of A B (functions-of A B N) = N
  using ⟨finite A⟩ ⟨finite B⟩ by (simp add: number-partition-of-functions-of)
  show number-partition-of A B ' ({f ∈ A →E B. f ' A = B} // domain-and-range-permutation A B)
  ⊆ {N. number-partition (card A) N ∧ size N = card B}
  using ⟨finite A⟩ ⟨finite B⟩ by (auto simp add: quotient-eq number-partition-of surj-on-implies-size-eq-card)
  show functions-of A B ' {N. number-partition (card A) N ∧ size N = card B}
  ⊆ {f ∈ A →E B. f ' A = B} // domain-and-range-permutation A B
  using ⟨finite A⟩ ⟨finite B⟩ by (auto simp add: quotient-eq intro: functions-of functions-of-is-surj-on)
qed

```

lemma *bij-betw-functions-of*:

```

  assumes finite A finite B
  shows bij-betw (functions-of A B) {N. number-partition (card A) N ∧ size N = card B} ({f ∈ A →E B. f ' A = B} // domain-and-range-permutation A B)
proof (rule bij-betw-byWitness[where f'=number-partition-of A B])
  have quotient-eq: {f ∈ A →E B. f ' A = B} // domain-and-range-permutation A B = {F ∈ ((A →E B) // domain-and-range-permutation A B). univ (λf. f ' A = B) F}
  using equiv-domain-and-range-permutation[of A B] surjective-respects-domain-and-range-permutation[of A B] by (simp only: univ-preserves-predicate)
  show ∀ F ∈ {f ∈ A →E B. f ' A = B} // domain-and-range-permutation A B.

```

```

    functions-of A B (number-partition-of A B F) = F
  using ⟨finite A⟩ ⟨finite B⟩ by (auto simp only: quotient-eq functions-of-number-partition-of)
  show  $\forall N \in \{N. \text{number-partition} (\text{card } A) N \wedge \text{size } N = \text{card } B\}. \text{number-partition-of}$ 
    A B (functions-of A B N) = N
  using ⟨finite A⟩ ⟨finite B⟩ by (simp add: number-partition-of-functions-of)
  show number-partition-of A B '  $\{f \in A \rightarrow_E B. f ' A = B\}$  // domain-and-range-permutation
    A B)
     $\subseteq \{N. \text{number-partition} (\text{card } A) N \wedge \text{size } N = \text{card } B\}$ 
  using ⟨finite A⟩ ⟨finite B⟩ by (auto simp add: quotient-eq number-partition-of
    surj-on-implies-size-eq-card)
  show functions-of A B '  $\{N. \text{number-partition} (\text{card } A) N \wedge \text{size } N = \text{card } B\}$ 
     $\subseteq \{f \in A \rightarrow_E B. f ' A = B\}$  // domain-and-range-permutation A B
  using ⟨finite A⟩ ⟨finite B⟩ by (auto simp add: quotient-eq intro: functions-of
    functions-of-is-surj-on)
qed

```

15.3 Cardinality

lemma *card-surjective-functions-domain-and-range-permutation:*

```

  assumes finite A finite B
  shows card  $\{f \in A \rightarrow_E B. f ' A = B\}$  // domain-and-range-permutation A B)
    = Partition (card A) (card B)
  proof -
    have bij-betw (number-partition-of A B)  $\{f \in A \rightarrow_E B. f ' A = B\}$  // do-
      main-and-range-permutation A B)  $\{N. \text{number-partition} (\text{card } A) N \wedge \text{size } N =$ 
        card B}
    using ⟨finite A⟩ ⟨finite B⟩ by (rule bij-betw-number-partition-of)
    from this have card  $\{f \in A \rightarrow_E B. f ' A = B\}$  // domain-and-range-permutation
      A B) = card  $\{N. \text{number-partition} (\text{card } A) N \wedge \text{size } N = \text{card } B\}$ 
    by (rule bij-betw-same-card)
    also have card  $\{N. \text{number-partition} (\text{card } A) N \wedge \text{size } N = \text{card } B\}$  = Partition
      (card A) (card B)
    by (rule card-partitions-with-k-parts)
    finally show ?thesis .
  qed

```

end

16 Cardinality of Bijections

theory *Card-Bijections*

imports

```

  Twelfold-Way-Entry2
  Twelfold-Way-Entry3
  Twelfold-Way-Entry5
  Twelfold-Way-Entry6
  Twelfold-Way-Entry8
  Twelfold-Way-Entry9
  Twelfold-Way-Entry11

```


begin

16.1 Bijections from A to B

lemma *bij-betw-set-is-empty:*

assumes *finite A finite B*

assumes *card A \neq card B*

shows $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} = \{\}$

using *assms bij-betw-same-card* **by** *blast*

lemma *card-bijections-eq-zero:*

assumes *finite A finite B*

assumes *card A \neq card B*

shows *card* $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} = 0$

using *bij-betw-set-is-empty*[*OF assms*] **by** (*simp only: card.empty*)

Two alternative proofs for the cardinality of bijections up to a permutation on A.

lemma

assumes *finite A finite B*

assumes *card A = card B*

shows *card* $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} = \text{fact } (\text{card } B)$

proof –

have *card* $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} = \text{card } \{f \in A \rightarrow_E B. \text{inj-on } f \ A\}$

using $\langle \text{finite } B \rangle \langle \text{card } A = \text{card } B \rangle$ **by** (*metis bij-betw-implies-inj-on-and-card-eq*)

also have $\dots = \text{fact } (\text{card } B)$

using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle \langle \text{card } A = \text{card } B \rangle$ **by** (*simp add: card-extensional-funcset-inj-on*)

finally show *?thesis* .

qed

lemma *card-bijections:*

assumes *finite A finite B*

assumes *card A = card B*

shows *card* $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} = \text{fact } (\text{card } B)$

proof –

have *card* $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} = \text{card } \{f \in A \rightarrow_E B. f \ A = B\}$

using $\langle \text{finite } A \rangle \langle \text{card } A = \text{card } B \rangle$

by (*metis bij-betw-implies-surj-on-and-card-eq*)

also have $\dots = \text{fact } (\text{card } B)$

using $\langle \text{finite } A \rangle \langle \text{finite } B \rangle \langle \text{card } A = \text{card } B \rangle$

by (*simp add: card-extensional-funcset-surj-on*)

finally show *?thesis* .

qed

16.2 Bijections from A to B up to a Permutation on A

lemma *bij-betw-quotient-domain-permutation-eq-empty:*

assumes *card A \neq card B*

shows $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B = \{\}$
using $\langle \text{card } A \neq \text{card } B \rangle \text{bij-betw-same-card}$ **by** *auto*

lemma *card-bijections-domain-permutation-eq-0:*

assumes $\text{card } A \neq \text{card } B$
shows $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B) = 0$
using *bij-betw-quotient-domain-permutation-eq-empty*[*OF assms*] **by** (*simp only: card.empty*)

Two alternative proofs for the cardinality of bijections up to a permutation on A.

lemma

assumes *finite A finite B*
assumes $\text{card } A = \text{card } B$
shows $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B) = 1$
proof –
from *assms* **have** $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B$
 $= \{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-permutation } A \ B$
by (*metis (no-types, lifting) PiE-cong bij-betw-implies-inj-on-and-card-eq*)
from this show *?thesis*
using *assms* **by** (*simp add: card-injective-functions-domain-permutation*)
qed

lemma *card-bijections-domain-permutation-eq-1:*

assumes *finite A finite B*
assumes $\text{card } A = \text{card } B$
shows $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B) = 1$
proof –
from *assms* **have** $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B$
 $= \{f \in A \rightarrow_E B. f \ A = B\} // \text{domain-permutation } A \ B$
by (*metis (no-types, lifting) PiE-cong bij-betw-implies-surj-on-and-card-eq*)
from this show *?thesis*
using *assms* **by** (*simp add: card-surjective-functions-domain-permutation*)
qed

lemma *card-bijections-domain-permutation:*

assumes *finite A finite B*
shows $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B) =$
iverson (card A = card B)
using *assms card-bijections-domain-permutation-eq-0 card-bijections-domain-permutation-eq-1*
unfolding *iverson-def* **by** *auto*

16.3 Bijections from A to B up to a Permutation on B

lemma *bij-betw-quotient-range-permutation-eq-empty:*

assumes $\text{card } A \neq \text{card } B$
shows $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B = \{\}$
using $\langle \text{card } A \neq \text{card } B \rangle \text{bij-betw-same-card}$ **by** *auto*

lemma *card-bijections-range-permutation-eq-0*:
assumes $\text{card } A \neq \text{card } B$
shows $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B) = 0$
using *bij-betw-quotient-range-permutation-eq-empty*[*OF assms*] **by** (*simp only: card.empty*)

Two alternative proofs for the cardinality of bijections up to a permutation on B.

lemma
assumes *finite A finite B*
assumes $\text{card } A = \text{card } B$
shows $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B) = 1$
proof –
from *assms* **have** $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B =$
 $\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{range-permutation } A \ B$
by (*metis (no-types, lifting) PiE-cong bij-betw-implies-inj-on-and-card-eq*)
from this show *?thesis*
using *assms* **by** (*simp add: iverson-def card-injective-functions-range-permutation*)
qed

lemma *card-bijections-range-permutation-eq-1*:
assumes *finite A finite B*
assumes $\text{card } A = \text{card } B$
shows $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B) = 1$
proof –
from *assms* **have** $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B =$
 $\{f \in A \rightarrow_E B. f \ ' \ A = B\} // \text{range-permutation } A \ B$
by (*metis (no-types, lifting) PiE-cong bij-betw-implies-surj-on-and-card-eq*)
from this show *?thesis*
using *assms* **by** (*simp add: card-surjective-functions-range-permutation*)
qed

lemma *card-bijections-range-permutation*:
assumes *finite A finite B*
shows $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B) = \text{iverson}$
 $(\text{card } A = \text{card } B)$
using *assms card-bijections-range-permutation-eq-0 card-bijections-range-permutation-eq-1*
unfolding *iverson-def* **by** *auto*

16.4 Bijections from A to B up to a Permutation on A and B

lemma *bij-betw-quotient-domain-and-range-permutation-eq-empty*:
assumes $\text{card } A \neq \text{card } B$
shows $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-and-range-permutation } A \ B = \{\}$
using $\langle \text{card } A \neq \text{card } B \rangle$ *bij-betw-same-card* **by** *auto*

lemma *card-bijections-domain-and-range-permutation-eq-0*:
assumes $\text{card } A \neq \text{card } B$

```

shows  $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-and-range-permutation } A \ B) = 0$ 
using bij-betw-quotient-domain-and-range-permutation-eq-empty[OF assms] by (simp only: card.empty)

```

Two alternative proofs for the cardinality of bijections up to a permutation on A and B.

```

lemma
  assumes finite A finite B
  assumes  $\text{card } A = \text{card } B$ 
  shows  $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-and-range-permutation } A \ B) = 1$ 
proof –
  from assms have  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-and-range-permutation } A \ B =$ 
     $\{f \in A \rightarrow_E B. \text{inj-on } f \ A\} // \text{domain-and-range-permutation } A \ B$ 
    by (metis (no-types, lifting) PiE-cong bij-betw-implies-inj-on-and-card-eq)
  from this show ?thesis
  using assms by (simp add: iverson-def card-injective-functions-domain-and-range-permutation)
qed

```

```

lemma card-bijections-domain-and-range-permutation-eq-1:
  assumes finite A finite B
  assumes  $\text{card } A = \text{card } B$ 
  shows  $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-and-range-permutation } A \ B) = 1$ 
proof –
  from assms have  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-and-range-permutation } A \ B =$ 
     $\{f \in A \rightarrow_E B. f \ A = B\} // \text{domain-and-range-permutation } A \ B$ 
    by (metis (no-types, lifting) PiE-cong bij-betw-implies-surj-on-and-card-eq)
  from this show ?thesis
  using assms by (simp add: card-surjective-functions-domain-and-range-permutation Partition-diag)
qed

```

```

lemma card-bijections-domain-and-range-permutation:
  assumes finite A finite B
  shows  $\text{card } (\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-and-range-permutation } A \ B) = \text{iverson } (\text{card } A = \text{card } B)$ 
using assms card-bijections-domain-and-range-permutation-eq-0 card-bijections-domain-and-range-permutation
unfolding iverson-def by auto

```

end

17 Direct Proofs for Cardinality of Bijections

```

theory Card-Bijections-Direct
imports

```

begin

17.1 Bijections from A to B up to a Permutation on A

17.1.1 Equivalence Class

lemma *bijections-in-domain-permutation:*

assumes *finite A finite B*

assumes *card A = card B*

shows $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} //$
domain-permutation A B

proof $-$

from *assms* **obtain** *f* **where** $f: f \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$

by (*metis finite-same-card-bij-on-ext-funcset mem-Collect-eq*)

moreover **have** *proj-f*: $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} = \text{domain-permutation } A \ B \text{ ``}\{f\}$

proof

from *f* **show** $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} \subseteq \text{domain-permutation } A \ B \text{ ``}\{f\}$

unfolding *domain-permutation-def*

by (*auto elim: obtain-domain-permutation-for-two-bijections*)

next

show *domain-permutation A B `` {f} $\subseteq \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$*

proof

fix *f'*

assume $f' \in \text{domain-permutation } A \ B \text{ ``}\{f\}$

have $(f', f) \in \text{domain-permutation } A \ B$

using $\langle f' \in \text{domain-permutation } A \ B \text{ ``}\{f\} \rangle \text{equiv-domain-permutation[of } A \ B]$

by (*simp add: equiv-class-eq-iff*)

from *this* **obtain** *p* **where** $p \text{ permutes } A \ \forall x \in A. f' \ x = f \ (p \ x)$

unfolding *domain-permutation-def* **by** *auto*

from *this* **have** $\text{bij-betw } (f \circ p) \ A \ B$

using *bij-betw-comp-iff f permutes-imp-bij* **by** *fastforce*

from *this* **have** $\text{bij-betw } f' \ A \ B$

using $\langle \forall x \in A. f' \ x = f \ (p \ x) \rangle$

by (*metis (mono-tags, lifting) bij-betw-cong comp-apply*)

moreover **have** $f' \in A \rightarrow_E B$

using $\langle f' \in \text{domain-permutation } A \ B \text{ ``}\{f\} \rangle$

unfolding *domain-permutation-def* **by** *auto*

ultimately **show** $f' \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$ **by** *simp*

qed

qed

ultimately **show** *?thesis* **by** (*simp add: quotientI*)

qed

lemma *bij-betw-quotient-domain-permutation-eq:*

assumes *finite A finite B*

assumes *card A = card B*

```

  shows  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B = \{\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}\}$ 
proof
  show  $\{\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}\} \subseteq \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} //$ 
  domain-permutation  $A \ B$ 
  by (simp add: bijections-in-domain-permutation[OF assms])
next
  show  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B \subseteq \{\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}\}$ 
proof
  fix  $F$ 
  assume  $F\text{-in: } F \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} //$  domain-permutation  $A \ B$ 
  have  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B = \{F \in ((A \rightarrow_E B) // \text{domain-permutation } A \ B). \text{univ } (\lambda f. \text{bij-betw } f \ A \ B) \ F\}$ 
  using equiv-domain-permutation[of  $A \ B$ ] bij-betw-respects-domain-permutation[of  $A \ B$ ] by (simp only: univ-preserves-predicate)
  from  $F\text{-in this}$  have  $F \in (A \rightarrow_E B) //$  domain-permutation  $A \ B$ 
  and  $\text{univ } (\lambda f. \text{bij-betw } f \ A \ B) \ F$ 
  by blast+
  have  $F = \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$ 
proof
  have  $\forall f \in F. f \in A \rightarrow_E B$ 
  using  $\langle F \in (A \rightarrow_E B) // \text{domain-permutation } A \ B \rangle$ 
  by (metis ImageE equiv-class-eq-iff equiv-domain-permutation quotientE)
  moreover have  $\forall f \in F. \text{bij-betw } f \ A \ B$ 
  using univ-predicate-impl-forall[OF equiv-domain-permutation bij-betw-respects-domain-permutation]
  using  $\langle F \in (A \rightarrow_E B) // \text{domain-permutation } A \ B \rangle \langle \text{univ } (\lambda f. \text{bij-betw } f \ A \ B) \ F \rangle$ 
  by auto
  ultimately show  $F \subseteq \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$  by auto
next
  show  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} \subseteq F$ 
proof
  fix  $f'$ 
  assume  $f' \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$ 
  from this have  $f' \in A \rightarrow_E B$   $\text{bij-betw } f' \ A \ B$  by auto
  obtain  $f$  where  $f \in A \rightarrow_E B$  and  $F = \text{domain-permutation } A \ B \text{ “ } \{f\}$ 
  using  $\langle F \in (A \rightarrow_E B) // \text{domain-permutation } A \ B \rangle$  by (auto elim: quotientE)
  have  $\text{bij-betw } f \ A \ B$ 
  using univ-commute'[OF equiv-domain-permutation bij-betw-respects-domain-permutation]
  using  $\langle f \in A \rightarrow_E B \rangle \langle F = \text{domain-permutation } A \ B \text{ “ } \{f\} \rangle \langle \text{univ } (\lambda f. \text{bij-betw } f \ A \ B) \ F \rangle$ 
  by auto
  obtain  $p$  where  $p$  permutes  $A \ \forall x \in A. f \ x = f' \ (p \ x)$ 
  using obtain-domain-permutation-for-two-bijections
  using  $\langle \text{bij-betw } f \ A \ B \rangle \langle \text{bij-betw } f' \ A \ B \rangle$  by blast
  from this  $\langle f \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle$ 
  have  $(f, f') \in \text{domain-permutation } A \ B$ 

```

```

    unfolding domain-permutation-def by auto
  from this show  $f' \in F$ 
    using  $\langle F = \text{domain-permutation } A \ B \ \text{"}\{f\}\rangle$  by simp
qed
qed
from this show  $F \in \{\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}\}$  by simp
qed
qed

```

17.1.2 Cardinality

```

lemma
  assumes finite A finite B
  assumes card A = card B
  shows card  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{domain-permutation } A \ B = 1$ 
  using bij-betw-quotient-domain-permutation-eq[OF assms] by auto

```

17.2 Bijections from A to B up to a Permutation on B

17.2.1 Equivalence Class

```

lemma bijections-in-range-permutation:
  assumes finite A finite B
  assumes card A = card B
  shows  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} //$ 
  range-permutation A B
proof -
  from assms obtain f where  $f: f \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$ 
  by (metis finite-same-card-bij-on-ext-funcset mem-Collect-eq)
  moreover have  $\text{proj-}f: \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} = \text{range-permutation } A \ B \ \text{"}\{f\}$ 
  proof
    from f show  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} \subseteq \text{range-permutation } A \ B \ \text{"}\{f\}$ 
    unfolding range-permutation-def
    by (auto elim: obtain-range-permutation-for-two-bijections)
  next
    show  $\text{range-permutation } A \ B \ \text{"}\{f\} \subseteq \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$ 
    proof
      fix f'
      assume  $f' \in \text{range-permutation } A \ B \ \text{"}\{f\}$ 
      have  $(f', f) \in \text{range-permutation } A \ B$ 
      using  $\langle f' \in \text{range-permutation } A \ B \ \text{"}\{f\}\rangle$  equiv-range-permutation[of A B]
      by (simp add: equiv-class-eq-iff)
      from this obtain p where p permutes B  $\forall x \in A. f' \ x = p \ (f \ x)$ 
      unfolding range-permutation-def by auto
      from this have  $\text{bij-betw } (p \circ f) \ A \ B$ 
      using bij-betw-comp-iff f permutes-imp-bij by fastforce
      from this have  $\text{bij-betw } f' \ A \ B$ 
      using  $\langle \forall x \in A. f' \ x = p \ (f \ x) \rangle$ 
      by (metis (mono-tags, lifting) bij-betw-cong comp-apply)
    qed
  qed

```

```

    moreover have  $f' \in A \rightarrow_E B$ 
    using  $\langle f' \in \text{range-permutation } A \ B \ \langle \{f\} \rangle$ 
    unfolding range-permutation-def by auto
    ultimately show  $f' \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$  by simp
  qed
qed
ultimately show ?thesis by (simp add: quotientI)
qed

lemma bij-betw-quotient-range-permutation-eq:
  assumes finite A finite B
  assumes card A = card B
  shows  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B = \{\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}\}$ 
proof
  show  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} \subseteq \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B$ 
  by (simp add: bijections-in-range-permutation[OF assms])
next
  show  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B \subseteq \{\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}\}$ 
proof
  fix F
  assume F-in:  $F \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B$ 
  have  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} // \text{range-permutation } A \ B = \{F \in ((A \rightarrow_E B) // \text{range-permutation } A \ B). \text{univ } (\lambda f. \text{bij-betw } f \ A \ B) \ F\}$ 
  using equiv-range-permutation[of A B] bij-betw-respects-range-permutation[of A B] by (simp only: univ-preserves-predicate)
  from this F-in have  $F \in (A \rightarrow_E B) // \text{range-permutation } A \ B$ 
  and  $\text{univ } (\lambda f. \text{bij-betw } f \ A \ B) \ F$  by blast+
  have  $F = \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$ 
proof
  have  $\forall f \in F. f \in A \rightarrow_E B$ 
  using  $\langle F \in (A \rightarrow_E B) // \text{range-permutation } A \ B \rangle$ 
  by (metis ImageE equiv-class-eq-iff equiv-range-permutation quotientE)
  moreover have  $\forall f \in F. \text{bij-betw } f \ A \ B$ 
  using univ-predicate-impl-forall[OF equiv-range-permutation bij-betw-respects-range-permutation]
  using  $\langle F \in (A \rightarrow_E B) // \text{range-permutation } A \ B \rangle \langle \text{univ } (\lambda f. \text{bij-betw } f \ A \ B) \ F \rangle$ 
  by auto
  ultimately show  $F \subseteq \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$  by auto
next
  show  $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\} \subseteq F$ 
proof
  fix  $f'$ 
  assume  $f' \in \{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$ 
  from this have  $f' \in A \rightarrow_E B$  bij-betw  $f' \ A \ B$  by auto
  obtain f where  $f \in A \rightarrow_E B$  and  $F = \text{range-permutation } A \ B \ \langle \{f\} \rangle$ 
  using  $\langle F \in (A \rightarrow_E B) // \text{range-permutation } A \ B \rangle$  by (auto elim: quotientE)

```



```

    have bij-betw f A B
    using univ-commute'[OF equiv-range-permutation bij-betw-respects-range-permutation]
      using  $\langle f \in A \rightarrow_E B \rangle \langle F = \text{range-permutation } A \ B \ \text{"}\{f\}\rangle \langle \text{univ } (\lambda f.$ 
bij-betw f A B) F
      by auto
    obtain p where p permutes B  $\forall x \in A. f\ x = p\ (f'\ x)$ 
    using obtain-range-permutation-for-two-bijections
    using  $\langle \text{bij-betw } f\ A\ B \rangle \langle \text{bij-betw } f'\ A\ B \rangle$  by blast
    from this  $\langle f \in A \rightarrow_E B \rangle \langle f' \in A \rightarrow_E B \rangle$ 
    have  $(f, f') \in \text{range-permutation } A\ B$ 
    unfolding range-permutation-def by auto
    from this show  $f' \in F$ 
    using  $\langle F = \text{range-permutation } A\ B \ \text{"}\{f\}\rangle$  by simp
  qed
qed
from this show  $F \in \{\{f \in A \rightarrow_E B. \text{bij-betw } f\ A\ B\}\}$  by simp
qed
qed

```

17.2.2 Cardinality

```

lemma card-bijections-range-permutation-eq-1:
  assumes finite A finite B
  assumes card A = card B
  shows card  $\{f \in A \rightarrow_E B. \text{bij-betw } f\ A\ B\} // \text{range-permutation } A\ B = 1$ 
  using bij-betw-quotient-range-permutation-eq[OF assms] by auto

```

17.3 Bijections from A to B up to a Permutation on A and B

17.3.1 Equivalence Class

```

lemma bijections-in-domain-and-range-permutation:
  assumes finite A finite B
  assumes card A = card B
  shows  $\{f \in A \rightarrow_E B. \text{bij-betw } f\ A\ B\} \in \{f \in A \rightarrow_E B. \text{bij-betw } f\ A\ B\} //$ 
domain-and-range-permutation A B
proof -
  from assms obtain f where f:  $f \in \{f \in A \rightarrow_E B. \text{bij-betw } f\ A\ B\}$ 
  by (metis finite-same-card-bij-on-ext-funcset mem-Collect-eq)
  moreover have proj-f:  $\{f \in A \rightarrow_E B. \text{bij-betw } f\ A\ B\} = \text{domain-and-range-permutation}$ 
A B  $\text{"}\{f\}$ 
  proof
    have id permutes A by (simp add: permutes-id)
    from f this show  $\{f \in A \rightarrow_E B. \text{bij-betw } f\ A\ B\} \subseteq \text{domain-and-range-permutation}$ 
A B  $\text{"}\{f\}$ 
    unfolding domain-and-range-permutation-def
    by (fastforce elim: obtain-range-permutation-for-two-bijections)
  next

```

```

  show domain-and-range-permutation A B “ {f} ⊆ {f ∈ A →E B. bij-betw f A
B}
  proof
    fix f'
    assume f' ∈ domain-and-range-permutation A B “ {f}
    have (f', f) ∈ domain-and-range-permutation A B
    using ⟨f' ∈ domain-and-range-permutation A B “ {f}⟩ equiv-domain-and-range-permutation[of
A B]
      by (simp add: equiv-class-eq-iff)
    from this obtain pA pB where pA permutes A pB permutes B
    and ∀ x ∈ A. f' x = pB (f (pA x))
    unfolding domain-and-range-permutation-def by auto
    from this have bij-betw (pB ∘ f ∘ pA) A B
    using bij-betw-comp-iff f permutes-imp-bij
    by (metis (no-types, lifting) mem-Collect-eq)
    from this have bij-betw f' A B
    using ⟨∀ x ∈ A. f' x = pB (f (pA x))⟩
    by (auto intro: bij-betw-congI)
    moreover have f' ∈ A →E B
    using ⟨f' ∈ domain-and-range-permutation A B “ {f}⟩
    unfolding domain-and-range-permutation-def by auto
    ultimately show f' ∈ {f ∈ A →E B. bij-betw f A B} by simp
  qed
qed
ultimately show ?thesis by (simp add: quotientI)
qed

lemma bij-betw-quotient-domain-and-range-permutation-eq:
  assumes finite A finite B
  assumes card A = card B
  shows {f ∈ A →E B. bij-betw f A B} // domain-and-range-permutation A B =
  {{f ∈ A →E B. bij-betw f A B}}
  proof
    show {{f ∈ A →E B. bij-betw f A B}}
      ⊆ {f ∈ A →E B. bij-betw f A B} // domain-and-range-permutation A B
    using bijections-in-domain-and-range-permutation[OF assms] by auto
  next
    show {f ∈ A →E B. bij-betw f A B} // domain-and-range-permutation A B ⊆
    {{f ∈ A →E B. bij-betw f A B}}
    proof
      fix F
      assume F-in: F ∈ {f ∈ A →E B. bij-betw f A B} // domain-and-range-permutation
A B
      have {f ∈ A →E B. bij-betw f A B} // domain-and-range-permutation A B =
      {F ∈ ((A →E B) // domain-and-range-permutation A B). univ (λf. bij-betw f A
B) F}
      using equiv-domain-and-range-permutation[of A B] bij-betw-respects-domain-and-range-permutation[of
A B] by (simp only: univ-preserves-predicate)
      from F-in this have F ∈ (A →E B) // domain-and-range-permutation A B

```

```

    and univ (λf. bij-betw f A B) F by blast+
  have F = {f ∈ A →E B. bij-betw f A B}
  proof
    have ∀f ∈ F. f ∈ A →E B
      using ⟨F ∈ (A →E B) // domain-and-range-permutation A B⟩
      by (metis ImageE equiv-class-eq-iff equiv-domain-and-range-permutation
quotientE)
    moreover have ∀f ∈ F. bij-betw f A B
      using univ-predicate-impl-forall[OF equiv-domain-and-range-permutation
bij-betw-respects-domain-and-range-permutation]
      using ⟨F ∈ (A →E B) // domain-and-range-permutation A B⟩ ⟨univ (λf.
bij-betw f A B) F⟩
      by auto
    ultimately show F ⊆ {f ∈ A →E B. bij-betw f A B} by auto
  next
    show {f ∈ A →E B. bij-betw f A B} ⊆ F
    proof
      fix f'
      assume f' ∈ {f ∈ A →E B. bij-betw f A B}
      from this have f' ∈ A →E B bij-betw f' A B by auto
      obtain f where f ∈ A →E B and F = domain-and-range-permutation A
B “ {f}
      using ⟨F ∈ (A →E B) // domain-and-range-permutation A B⟩ by (auto
elim: quotientE)
      have bij-betw f A B
      using univ-commute'[OF equiv-domain-and-range-permutation bij-betw-respects-domain-and-range-perm
      using ⟨f ∈ A →E B⟩ ⟨F = domain-and-range-permutation A B “ {f}⟩
      ⟨univ (λf. bij-betw f A B) F⟩
      by auto
      obtain p where p permutes A ∀x∈A. f x = f' (p x)
      using obtain-domain-permutation-for-two-bijections
      using ⟨bij-betw f A B⟩ ⟨bij-betw f' A B⟩ by blast
      moreover have id permutes B by (simp add: permutes-id)
      moreover note ⟨f ∈ A →E B⟩ ⟨f' ∈ A →E B⟩
      ultimately have (f, f') ∈ domain-and-range-permutation A B
      unfolding domain-and-range-permutation-def id-def by auto
      from this show f' ∈ F
      using ⟨F = domain-and-range-permutation A B “ {f}⟩ by simp
    qed
  qed
  from this show F ∈ {{f ∈ A →E B. bij-betw f A B}} by simp
qed
qed

```

17.3.2 Cardinality

lemma *card-bijections-domain-and-range-permutation-eq-1:*

assumes *finite A finite B*
 assumes *card A = card B*

```

shows card ( $\{f \in A \rightarrow_E B. \text{bij-betw } f \ A \ B\}$  // domain-and-range-permutation
 $A \ B$ ) = 1
using bij-betw-quotient-domain-and-range-permutation-eq[OF assms] by auto

end

```

18 The Twelfold Way

```

theory Twelfold-Way
imports
  Preliminaries
  Twelfold-Way-Core
  Equiv-Relations-on-Functions
  Twelfold-Way-Entry1
  Twelfold-Way-Entry2
  Twelfold-Way-Entry4
  Twelfold-Way-Entry5
  Twelfold-Way-Entry6
  Twelfold-Way-Entry7
  Twelfold-Way-Entry8
  Twelfold-Way-Entry9
  Twelfold-Way-Entry3
  Twelfold-Way-Entry10
  Twelfold-Way-Entry11
  Twelfold-Way-Entry12
  Card-Bijections
  Card-Bijections-Direct
begin

end

```

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