

# Taylor Models

Christoph Traut and Fabian Immler

June 16, 2019

## Abstract

We present a formally verified implementation of multivariate Taylor models. Taylor models are a form of rigorous polynomial approximation, consisting of an approximation polynomial based on Taylor expansions, combined with a rigorous bound on the approximation error. Taylor models were introduced as a tool to mitigate the dependency problem of interval arithmetic. Our implementation automatically computes Taylor models for the class of elementary functions, expressed by composition of arithmetic operations and basic functions like exp, sin, or square root.

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## 1 Topology for Floating Point Numbers

**theory** *Float-Topology*

**imports**

*HOL-Analysis.Analysis*

*HOL-Library.Float*

**begin**

This topology is totally disconnected and not complete, in which sense is it useful? Perhaps for convergence of intervals?

**unbundle** *float.lifting*

**instantiation** *float :: dist*

**begin**

**lift-definition** *dist-float :: float  $\Rightarrow$  float  $\Rightarrow$  real is dist* *<proof>*

**lemma** *dist-float-eq-0-iff: (dist x y = 0) = (x = y) for x y::float*  
*<proof>*

**lemma** *dist-float-triangle2: dist x y  $\leq$  dist x z + dist y z for x y z::float*  
*<proof>*

**instance** *<proof>*

**end**

**instantiation** *float :: uniformity*

**begin**

**definition** *uniformity-float :: (float  $\times$  float) filter*

**where** *uniformity-float = (INF e $\in$ {0<..}. principal {(x, y). dist x y < e})*

**instance** *<proof>*

**end**

**lemma** *float-dense-in-real:*

**fixes** *x :: real*

**assumes** *x < y*

**shows**  *$\exists r \in \text{float}. x < r \wedge r < y$*

*<proof>*

```

lemma real-of-float-dense:
  fixes  $x\ y :: \text{real}$ 
  assumes  $x < y$ 
  shows  $\exists q :: \text{float}. x < \text{real-of-float } q \wedge \text{real-of-float } q < y$ 
   $\langle \text{proof} \rangle$ 

instantiation float :: linorder-topology
begin

definition open-float::float set  $\Rightarrow$  bool where
  open-float  $S = (\forall x \in S. \exists e > 0. \forall y. \text{dist } y\ x < e \longrightarrow y \in S)$ 

instance
 $\langle \text{proof} \rangle$ 

end

instance float :: metric-space
 $\langle \text{proof} \rangle$ 

instance float::topological-ab-group-add
 $\langle \text{proof} \rangle$ 

lifting-update float.lifting
lifting-forget float.lifting

end

```

## 2 Interval Type

```

theory Interval
  imports
    HOL-Analysis.Analysis
    HOL-Library.Set-Algebras
    HOL-Library.Float
  begin

  A type of non-empty, closed intervals.

  typedef (overloaded) 'a interval =
     $\{(a::'a::\text{preorder}, b). a \leq b\}$ 
    morphisms bounds-of-interval Interval
     $\langle \text{proof} \rangle$ 

  setup-lifting type-definition-interval

  lift-definition lower::('a::preorder) interval  $\Rightarrow$  'a is fst  $\langle \text{proof} \rangle$ 

  lift-definition upper::('a::preorder) interval  $\Rightarrow$  'a is snd  $\langle \text{proof} \rangle$ 

```

**lemma** *interval-eq-iff*:  $a = b \iff \text{lower } a = \text{lower } b \wedge \text{upper } a = \text{upper } b$   
*<proof>*

**lemma** *interval-eqI*:  $\text{lower } a = \text{lower } b \implies \text{upper } a = \text{upper } b \implies a = b$   
*<proof>*

**lemma** *lower-le-upper[simp]*:  $\text{lower } i \leq \text{upper } i$   
*<proof>*

**lift-definition** *set-of* ::  $'a::\text{preorder interval} \Rightarrow 'a \text{ set}$  **is**  $\lambda x. \{\text{fst } x \dots \text{snd } x\}$  *<proof>*

**lemma** *set-of-eq*:  $\text{set-of } x = \{\text{lower } x \dots \text{upper } x\}$   
*<proof>*

**context notes** *[[typedef-overloaded]]* **begin**

**lift-definition**(*code-dt*) *Interval* ::  $'a::\text{preorder} \Rightarrow 'a::\text{preorder} \Rightarrow 'a \text{ interval option}$   
**is**  $\lambda a b. \text{if } a \leq b \text{ then Some } (a, b) \text{ else None}$   
*<proof>*

**end**

**instantiation** *interval* ::  $(\{\text{preorder}, \text{equal}\}) \text{ equal}$   
**begin**

**definition** *equal-class.equal*  $a b \equiv (\text{lower } a = \text{lower } b) \wedge (\text{upper } a = \text{upper } b)$

**instance** *<proof>*  
**end**

**instantiation** *interval* ::  $(\text{preorder}) \text{ ord}$  **begin**

**definition** *less-eq-interval* ::  $'a \text{ interval} \Rightarrow 'a \text{ interval} \Rightarrow \text{bool}$   
**where**  $\text{less-eq-interval } a b \iff \text{lower } b \leq \text{lower } a \wedge \text{upper } a \leq \text{upper } b$

**definition** *less-interval* ::  $'a \text{ interval} \Rightarrow 'a \text{ interval} \Rightarrow \text{bool}$   
**where**  $\text{less-interval } x y = (x \leq y \wedge \neg y \leq x)$

**instance** *<proof>*  
**end**

**instantiation** *interval* ::  $(\text{lattice}) \text{ semilattice-sup}$   
**begin**

**lift-definition** *sup-interval* ::  $'a \text{ interval} \Rightarrow 'a \text{ interval} \Rightarrow 'a \text{ interval}$   
**is**  $\lambda(a, b) (c, d). (\text{inf } a c, \text{sup } b d)$   
*<proof>*

**lemma** *lower-sup[simp]*:  $\text{lower } (\text{sup } A B) = \text{inf } (\text{lower } A) (\text{lower } B)$

$\langle proof \rangle$

**lemma** *upper-sup[simp]: upper (sup A B) = sup (upper A) (upper B)*  
 $\langle proof \rangle$

**instance**  $\langle proof \rangle$   
**end**

**lemma** *set-of-interval-union: set-of A  $\cup$  set-of B  $\subseteq$  set-of (sup A B) for A::'a::lattice interval*  
 $\langle proof \rangle$

**lemma** *interval-union-commute: sup A B = sup B A for A::'a::lattice interval*  
 $\langle proof \rangle$

**lemma** *interval-union-mono1: set-of a  $\subseteq$  set-of (sup a A) for A :: 'a::lattice interval*  
 $\langle proof \rangle$

**lemma** *interval-union-mono2: set-of A  $\subseteq$  set-of (sup a A) for A :: 'a::lattice interval*  
 $\langle proof \rangle$

**lift-definition** *interval-of :: 'a::preorder  $\Rightarrow$  'a interval is  $\lambda x. (x, x)$*   
 $\langle proof \rangle$

**lemma** *lower-interval-of[simp]: lower (interval-of a) = a*  
 $\langle proof \rangle$

**lemma** *upper-interval-of[simp]: upper (interval-of a) = a*  
 $\langle proof \rangle$

**definition** *width :: 'a::{preorder,minus} interval  $\Rightarrow$  'a*  
**where** *width i = upper i - lower i*

**instantiation** *interval :: (ordered-ab-semigroup-add) ab-semigroup-add*  
**begin**

**lift-definition** *plus-interval::'a interval  $\Rightarrow$  'a interval  $\Rightarrow$  'a interval*  
**is**  $\lambda(a, b). \lambda(c, d). (a + c, b + d)$   
 $\langle proof \rangle$

**lemma** *lower-plus[simp]: lower (plus A B) = plus (lower A) (lower B)*  
 $\langle proof \rangle$

**lemma** *upper-plus[simp]: upper (plus A B) = plus (upper A) (upper B)*  
 $\langle proof \rangle$

**instance**  $\langle proof \rangle$   
**end**

**instance** *interval* :: (*{ordered-ab-semigroup-add, lattice}*) *ordered-ab-semigroup-add*  
⟨*proof*⟩

**instantiation** *interval* :: (*{preorder, zero}*) *zero*  
**begin**

**lift-definition** *zero-interval*::'a *interval* **is** (*0, 0*) ⟨*proof*⟩

**lemma** *lower-zero*[*simp*]: *lower 0 = 0*  
⟨*proof*⟩

**lemma** *upper-zero*[*simp*]: *upper 0 = 0*  
⟨*proof*⟩

**instance** ⟨*proof*⟩  
**end**

**instance** *interval* :: (*{ordered-comm-monoid-add}*) *comm-monoid-add*  
⟨*proof*⟩

**instance** *interval* :: (*{ordered-comm-monoid-add, lattice}*) *ordered-comm-monoid-add*  
⟨*proof*⟩

**instantiation** *interval* :: (*{ordered-ab-group-add}*) *uminus*  
**begin**

**lift-definition** *uminus-interval*::'a *interval*  $\Rightarrow$  'a *interval* **is**  $\lambda(a, b). (-b, -a)$   
⟨*proof*⟩

**lemma** *lower-uminus*[*simp*]: *lower (- A) = - upper A*  
⟨*proof*⟩

**lemma** *upper-uminus*[*simp*]: *upper (- A) = - lower A*  
⟨*proof*⟩

**instance** ⟨*proof*⟩  
**end**

**instantiation** *interval* :: (*{ordered-ab-group-add}*) *minus*  
**begin**

**definition** *minus-interval*::'a *interval*  $\Rightarrow$  'a *interval*  $\Rightarrow$  'a *interval*

**where** *minus-interval a b = a + - b*

**lemma** *lower-minus*[*simp*]: *lower (minus A B) = minus (lower A) (upper B)*  
⟨*proof*⟩

**lemma** *upper-minus*[*simp*]: *upper (minus A B) = minus (upper A) (lower B)*  
⟨*proof*⟩

**instance** ⟨*proof*⟩  
**end**

**instantiation** *interval* :: (*linordered-semiring*) *times*  
**begin**

**lift-definition** *times-interval* :: 'a interval  $\Rightarrow$  'a interval  $\Rightarrow$  'a interval  
**is**  $\lambda(a1, a2). \lambda(b1, b2).$   
 (let  $x1 = a1 * b1; x2 = a1 * b2; x3 = a2 * b1; x4 = a2 * b2$   
 in (min  $x1$  (min  $x2$  (min  $x3$   $x4$ )), max  $x1$  (max  $x2$  (max  $x3$   $x4$ ))))  
 <proof>

**lemma** *lower-times*:

$lower (times A B) = Min \{lower A * lower B, lower A * upper B, upper A * lower B, upper A * upper B\}$   
 <proof>

**lemma** *upper-times*:

$upper (times A B) = Max \{lower A * lower B, lower A * upper B, upper A * lower B, upper A * upper B\}$   
 <proof>

**instance** <proof>

**end**

**lemma** *interval-eq-set-of-iff*:  $X = Y \longleftrightarrow set-of X = set-of Y$  **for**  $X Y :: 'a :: order interval$   
 <proof>

## 2.1 Membership

**abbreviation** (in *preorder*) *in-interval* ((-/  $\in_i$  -) [51, 51] 50)  
 where *in-interval*  $x X \equiv x \in set-of X$

**lemma** *in-interval-to-interval*[intro!]:  $a \in_i interval-of a$   
 <proof>

**lemma** *plus-in-intervalI*:

**fixes**  $x y :: 'a :: ordered-ab-semigroup-add$   
**shows**  $x \in_i X \Longrightarrow y \in_i Y \Longrightarrow x + y \in_i X + Y$   
 <proof>

**lemma** *connected-set-of*[intro, simp]:

$connected (set-of X)$  **for**  $X :: 'a :: linear-continuum-topology interval$   
 <proof>

**lemma** *ex-sum-in-interval-lemma*:  $\exists xa \in \{la .. ua\}. \exists xb \in \{lb .. ub\}. x = xa + xb$

**if**  $la \leq ua \ lb \leq ub \ la + lb \leq x \ x \leq ua + ub$

$ua - la \leq ub - lb$

**for**  $la \ b \ c \ d :: 'a :: linordered-ab-group-add$

<proof>

**lemma** *ex-sum-in-interval*:  $\exists xa \geq la. xa \leq ua \wedge (\exists xb \geq lb. xb \leq ub \wedge x = xa + xb)$

**if**  $a: la \leq ua$  **and**  $b: lb \leq ub$  **and**  $x: la + lb \leq x \leq ua + ub$

**for**  $la\ b\ c\ d :: 'a :: \text{linordered-ab-group-add}$   
 $\langle \text{proof} \rangle$

**lemma** *Icc-plus-Icc*:

$\{a .. b\} + \{c .. d\} = \{a + c .. b + d\}$   
**if**  $a \leq b\ c \leq d$   
**for**  $a\ b\ c\ d :: 'a :: \text{linordered-ab-group-add}$   
 $\langle \text{proof} \rangle$

**lemma** *set-of-plus*:

**fixes**  $A :: 'a :: \text{linordered-ab-group-add interval}$   
**shows**  $\text{set-of } (A + B) = \text{set-of } A + \text{set-of } B$   
 $\langle \text{proof} \rangle$

**lemma** *plus-in-intervalE*:

**fixes**  $xy :: 'a :: \text{linordered-ab-group-add}$   
**assumes**  $xy \in_i X + Y$   
**obtains**  $x\ y$  **where**  $xy = x + y\ x \in_i X\ y \in_i Y$   
 $\langle \text{proof} \rangle$

**lemma** *set-of-uminus*:  $\text{set-of } (-X) = \{-x \mid x. x \in \text{set-of } X\}$

**for**  $X :: 'a :: \text{ordered-ab-group-add interval}$   
 $\langle \text{proof} \rangle$

**lemma** *uminus-in-intervalI*:

**fixes**  $x :: 'a :: \text{ordered-ab-group-add}$   
**shows**  $x \in_i X \implies -x \in_i -X$   
 $\langle \text{proof} \rangle$

**lemma** *uminus-in-intervalD*:

**fixes**  $x :: 'a :: \text{ordered-ab-group-add}$   
**shows**  $x \in_i -X \implies -x \in_i X$   
 $\langle \text{proof} \rangle$

**lemma** *minus-in-intervalI*:

**fixes**  $x\ y :: 'a :: \text{ordered-ab-group-add}$   
**shows**  $x \in_i X \implies y \in_i Y \implies x - y \in_i X - Y$   
 $\langle \text{proof} \rangle$

**lemma** *set-of-minus*:  $\text{set-of } (X - Y) = \{x - y \mid x\ y. x \in \text{set-of } X \wedge y \in \text{set-of } Y\}$

**for**  $X\ Y :: 'a :: \text{linordered-ab-group-add interval}$   
 $\langle \text{proof} \rangle$

**lemma** *times-in-intervalI*:

**fixes**  $x\ y :: 'a :: \text{linordered-ring}$   
**assumes**  $x \in_i X\ y \in_i Y$   
**shows**  $x * y \in_i X * Y$   
 $\langle \text{proof} \rangle$

**lemma** *times-in-intervalE*:  
**fixes**  $xy :: 'a :: \{\text{linordered-semiring, real-normed-algebra, linear-continuum-topology}\}$   
— TODO: linear continuum topology is pretty strong  
**assumes**  $xy \in_i X * Y$   
**obtains**  $x y$  **where**  $xy = x * y$   $x \in_i X$   $y \in_i Y$   
 $\langle \text{proof} \rangle$

**lemma** *set-of-times*:  $\text{set-of } (X * Y) = \{x * y \mid x y. x \in \text{set-of } X \wedge y \in \text{set-of } Y\}$   
**for**  $X Y :: 'a :: \{\text{linordered-ring, real-normed-algebra, linear-continuum-topology}\}$   
*interval*  
 $\langle \text{proof} \rangle$

**instance** *interval* :: (*linordered-idom*) *cancel-semigroup-add*  
 $\langle \text{proof} \rangle$

**lemma** *interval-mul-commute*:  $A * B = B * A$  **for**  $A B :: 'a :: \text{linordered-idom}$   
*interval*  
 $\langle \text{proof} \rangle$

**lemma** *interval-times-zero-right[simp]*:  $A * 0 = 0$  **for**  $A :: 'a :: \text{linordered-ring}$   
*interval*  
 $\langle \text{proof} \rangle$

**lemma** *interval-times-zero-left[simp]*:  
 $0 * A = 0$  **for**  $A :: 'a :: \text{linordered-ring}$  *interval*  
 $\langle \text{proof} \rangle$

**instantiation** *interval* :: ( $\{\text{preorder, one}\}$ ) *one*  
**begin**

**lift-definition** *one-interval*:: $'a$  *interval* **is**  $(1, 1)$   $\langle \text{proof} \rangle$   
**lemma** *lower-one[simp]*: *lower*  $1 = 1$   
 $\langle \text{proof} \rangle$   
**lemma** *upper-one[simp]*: *upper*  $1 = 1$   
 $\langle \text{proof} \rangle$   
**instance**  $\langle \text{proof} \rangle$   
**end**

**instance** *interval* :: ( $\{\text{one, preorder, linordered-semiring}\}$ ) *power*  
 $\langle \text{proof} \rangle$

**lemma** *set-of-one[simp]*:  $\text{set-of } (1 :: 'a :: \{\text{one, order}\} \text{ interval}) = \{1\}$   
 $\langle \text{proof} \rangle$

**instance** *interval* ::  
(*linordered-idom, linordered-ring, real-normed-algebra, linear-continuum-topology*)  
*monoid-mult*  
 $\langle \text{proof} \rangle$

**lemma** *one-times-ivl-left[simp]*:  $1 * A = A$  **for**  $A :: 'a::linordered-idom$  interval  
⟨proof⟩

**lemma** *one-times-ivl-right[simp]*:  $A * 1 = A$  **for**  $A :: 'a::linordered-idom$  interval  
⟨proof⟩

**lemma** *set-of-power-mono*:  $a^n \in \text{set-of } (A^n)$  **if**  $a \in \text{set-of } A$   
**for**  $a :: 'a::linordered-idom$   
⟨proof⟩

**lemma** *set-of-add-cong*:  
 $\text{set-of } (A + B) = \text{set-of } (A' + B')$   
**if**  $\text{set-of } A = \text{set-of } A'$   $\text{set-of } B = \text{set-of } B'$   
**for**  $A :: 'a::linordered-ab-group-add$  interval  
⟨proof⟩

**lemma** *set-of-add-inc-left*:  
 $\text{set-of } (A + B) \subseteq \text{set-of } (A' + B)$   
**if**  $\text{set-of } A \subseteq \text{set-of } A'$   
**for**  $A :: 'a::linordered-ab-group-add$  interval  
⟨proof⟩

**lemma** *set-of-add-inc-right*:  
 $\text{set-of } (A + B) \subseteq \text{set-of } (A + B')$   
**if**  $\text{set-of } B \subseteq \text{set-of } B'$   
**for**  $A :: 'a::linordered-ab-group-add$  interval  
⟨proof⟩

**lemma** *set-of-add-inc*:  
 $\text{set-of } (A + B) \subseteq \text{set-of } (A' + B')$   
**if**  $\text{set-of } A \subseteq \text{set-of } A'$   $\text{set-of } B \subseteq \text{set-of } B'$   
**for**  $A :: 'a::linordered-ab-group-add$  interval  
⟨proof⟩

**lemma** *set-of-neg-inc*:  
 $\text{set-of } (-A) \subseteq \text{set-of } (-A')$   
**if**  $\text{set-of } A \subseteq \text{set-of } A'$   
**for**  $A :: 'a::ordered-ab-group-add$  interval  
⟨proof⟩

**lemma** *set-of-sub-inc-left*:  
 $\text{set-of } (A - B) \subseteq \text{set-of } (A' - B)$   
**if**  $\text{set-of } A \subseteq \text{set-of } A'$   
**for**  $A :: 'a::linordered-ab-group-add$  interval  
⟨proof⟩

**lemma** *set-of-sub-inc-right*:  
 $\text{set-of } (A - B) \subseteq \text{set-of } (A - B')$

**if**  $set-of\ B \subseteq set-of\ B'$   
**for**  $A :: 'a::linordered-ab-group-add\ interval$   
 $\langle proof \rangle$

**lemma** *set-of-sub-inc:*  
 $set-of\ (A - B) \subseteq set-of\ (A' - B')$   
**if**  $set-of\ A \subseteq set-of\ A'\ set-of\ B \subseteq set-of\ B'$   
**for**  $A :: 'a::linordered-idom\ interval$   
 $\langle proof \rangle$

**lemma** *set-of-mul-inc-right:*  
 $set-of\ (A * B) \subseteq set-of\ (A * B')$   
**if**  $set-of\ B \subseteq set-of\ B'$   
**for**  $A :: 'a::linordered-ring\ interval$   
 $\langle proof \rangle$

**lemma** *set-of-distrib-left:*  
 $set-of\ (B * (A1 + A2)) \subseteq set-of\ (B * A1 + B * A2)$   
**for**  $A1 :: 'a::linordered-ring\ interval$   
 $\langle proof \rangle$

**lemma** *set-of-distrib-right:*  
 $set-of\ ((A1 + A2) * B) \subseteq set-of\ (A1 * B + A2 * B)$   
**for**  $A1\ A2\ B :: 'a::\{linordered-ring, real-normed-algebra, linear-continuum-topology\}$   
 $interval$   
 $\langle proof \rangle$

**lemma** *set-of-mul-inc-left:*  
 $set-of\ (A * B) \subseteq set-of\ (A' * B)$   
**if**  $set-of\ A \subseteq set-of\ A'$   
**for**  $A :: 'a::\{linordered-ring, real-normed-algebra, linear-continuum-topology\}\ interval$   
 $\langle proof \rangle$

**lemma** *set-of-mul-inc:*  
 $set-of\ (A * B) \subseteq set-of\ (A' * B')$   
**if**  $set-of\ A \subseteq set-of\ A'\ set-of\ B \subseteq set-of\ B'$   
**for**  $A :: 'a::\{linordered-ring, real-normed-algebra, linear-continuum-topology\}\ interval$   
 $\langle proof \rangle$

**lemma** *set-of-pow-inc:*  
 $set-of\ (A ^n) \subseteq set-of\ (A' ^n)$   
**if**  $set-of\ A \subseteq set-of\ A'$   
**for**  $A :: 'a::\{linordered-idom, real-normed-algebra, linear-continuum-topology\}$   
 $interval$   
 $\langle proof \rangle$

**lemma** *set-of-distrib-right-left:*  
 $set-of\ ((A1 + A2) * (B1 + B2)) \subseteq set-of\ (A1 * B1 + A1 * B2 + A2 * B1 + A2 * B2)$

**for**  $A1 :: 'a::\{linordered-idom, real-normed-algebra, linear-continuum-topology\}$   
*interval*  
 <proof>

**lemma** *mult-bounds-enclose-zero1*:  
 $\min (la * lb) (\min (la * ub) (\min (lb * ua) (ua * ub))) \leq 0$   
 $0 \leq \max (la * lb) (\max (la * ub) (\max (lb * ua) (ua * ub)))$   
**if**  $la \leq 0 \ 0 \leq ua$   
**for**  $la \ lb \ ua \ ub :: 'a::linordered-idom$   
 <proof>

**lemma** *mult-bounds-enclose-zero2*:  
 $\min (la * lb) (\min (la * ub) (\min (lb * ua) (ua * ub))) \leq 0$   
 $0 \leq \max (la * lb) (\max (la * ub) (\max (lb * ua) (ua * ub)))$   
**if**  $lb \leq 0 \ 0 \leq ub$   
**for**  $la \ lb \ ua \ ub :: 'a::linordered-idom$   
 <proof>

**lemma** *set-of-mul-contains-zero*:  
 $0 \in \text{set-of } (A * B)$   
**if**  $0 \in \text{set-of } A \vee 0 \in \text{set-of } B$   
**for**  $A :: 'a::linordered-idom \text{ interval}$   
 <proof>

**instance** *interval* :: *(linordered-semiring) mult-zero*  
 <proof>

**lift-definition** *min-interval*:: $'a::linorder \text{ interval} \Rightarrow 'a \text{ interval} \Rightarrow 'a \text{ interval}$  **is**  
 $\lambda(l1, u1). \lambda(l2, u2). (\min l1 \ l2, \min u1 \ u2)$   
 <proof>

**lemma** *lower-min-interval[simp]*:  $\text{lower } (\text{min-interval } x \ y) = \min (\text{lower } x) (\text{lower } y)$   
 <proof>

**lemma** *upper-min-interval[simp]*:  $\text{upper } (\text{min-interval } x \ y) = \min (\text{upper } x) (\text{upper } y)$   
 <proof>

**lemma** *min-intervalI*:  
 $a \in_i A \Longrightarrow b \in_i B \Longrightarrow \min a \ b \in_i \text{min-interval } A \ B$   
 <proof>

**lift-definition** *max-interval*:: $'a::linorder \text{ interval} \Rightarrow 'a \text{ interval} \Rightarrow 'a \text{ interval}$  **is**  
 $\lambda(l1, u1). \lambda(l2, u2). (\max l1 \ l2, \max u1 \ u2)$   
 <proof>

**lemma** *lower-max-interval[simp]*:  $\text{lower } (\text{max-interval } x \ y) = \max (\text{lower } x) (\text{lower } y)$   
 <proof>

**lemma** *upper-max-interval[simp]*:  $\text{upper } (\text{max-interval } x \ y) = \max (\text{upper } x) (\text{upper } y)$

*<proof>*

**lemma** *max-intervalI*:

$a \in_i A \implies b \in_i B \implies \max a b \in_i \max\text{-interval } A B$

*<proof>*

**lift-definition** *abs-interval*::*'a::linordered-idom interval*  $\Rightarrow$  *'a interval is*

$(\lambda(l,u). (\text{if } l < 0 \wedge 0 < u \text{ then } 0 \text{ else } \min |l| |u|, \max |l| |u|))$

*<proof>*

**lemma** *lower-abs-interval[simp]*:

$\text{lower } (\text{abs-interval } x) = (\text{if } \text{lower } x < 0 \wedge 0 < \text{upper } x \text{ then } 0 \text{ else } \min |\text{lower } x| |\text{upper } x|)$

*<proof>*

**lemma** *upper-abs-interval[simp]*:  $\text{upper } (\text{abs-interval } x) = \max |\text{lower } x| |\text{upper } x|$

*<proof>*

**lemma** *in-abs-intervalI1*:

$lx < 0 \implies 0 < ux \implies 0 \leq xa \implies xa \leq \max (-lx) (ux) \implies xa \in \text{abs } \{lx..ux\}$

**for**  $xa::'a::\text{linordered-idom}$

*<proof>*

**lemma** *in-abs-intervalI2*:

$\min (|lx|) |ux| \leq xa \implies xa \leq \max |lx| |ux| \implies lx \leq ux \implies 0 \leq lx \vee ux \leq 0$

$\implies$

$xa \in \text{abs } \{lx..ux\}$

**for**  $xa::'a::\text{linordered-idom}$

*<proof>*

**lemma** *set-of-abs-interval*:  $\text{set-of } (\text{abs-interval } x) = \text{abs } \{ \text{set-of } x$

*<proof>*

**fun** *split-domain* :: (*'a::preorder interval*  $\Rightarrow$  *'a interval list*)  $\Rightarrow$  *'a interval list*  $\Rightarrow$  *'a interval list list*

**where** *split-domain split* [] = [[]]

| *split-domain split* (I#Is) = (

  let S = *split* I;

  D = *split-domain split* Is

  in *concat* (*map* ( $\lambda d. \text{map } (\lambda s. s \# d) S$ ) D)

)

**context notes** [[*typedef-overloaded*]] **begin**

**lift-definition**(*code-dt*) *split-interval*::*'a::linorder interval*  $\Rightarrow$  *'a*  $\Rightarrow$  (*'a interval*  $\times$  *'a interval*)

**is**  $\lambda(l, u) x. ((\min l x, \max l x), (\min u x, \max u x))$

*<proof>*

**end**

**lemma** *split-domain-nonempty*:  
**assumes**  $\bigwedge I. \text{split } I \neq []$   
**shows** *split-domain split*  $I \neq []$   
 $\langle \text{proof} \rangle$

**lemma** *split-intervalD*: *split-interval*  $X \ x = (A, B) \implies \text{set-of } X \subseteq \text{set-of } A \cup \text{set-of } B$   
 $\langle \text{proof} \rangle$

**definition** *split-float-interval*  $x = \text{split-interval } x ((\text{lower } x + \text{upper } x) * \text{Float } 1 (-1))$

**lemma** *split-float-intervalD*: *split-float-interval*  $X = (A, B) \implies \text{set-of } X \subseteq \text{set-of } A \cup \text{set-of } B$   
 $\langle \text{proof} \rangle$

**lemmas** *float-round-down-le*[*intro*] = *order-trans*[*OF float-round-down*]  
**and** *float-round-up-ge*[*intro*] = *order-trans*[*OF - float-round-up*]

**instantiation** *interval* :: ( $\{\text{topological-space, preorder}\}$ ) *topological-space*  
**begin**

**definition** *open-interval-def*[*code del*]: *open* ( $X :: 'a \text{ interval set}$ ) =  
 $(\forall x \in X.$   
 $\quad \exists A \ B.$   
 $\quad \quad \text{open } A \wedge$   
 $\quad \quad \text{open } B \wedge$   
 $\quad \quad \text{lower } x \in A \wedge \text{upper } x \in B \wedge \text{Interval } ' (A \times B) \subseteq X)$

**instance**  
 $\langle \text{proof} \rangle$

**end**

**definition** *mid* :: *float interval*  $\Rightarrow$  *float*  
**where** *mid*  $i = (\text{lower } i + \text{upper } i) * \text{Float } 1 (-1)$

**lemma** *mid-in-interval*: *mid*  $i \in_i i$   
 $\langle \text{proof} \rangle$

**definition** *centered* :: *float interval*  $\Rightarrow$  *float interval*  
**where** *centered*  $i = i - \text{interval-of } (\text{mid } i)$

## 2.2 Quickcheck

**lift-definition** *Ivl*:: $'a \Rightarrow 'a :: \text{preorder} \Rightarrow 'a \text{ interval}$  **is**  $\lambda a \ b. (\text{min } a \ b, b)$   
 $\langle \text{proof} \rangle$

**instantiation** *interval* :: (*{exhaustive,preorder}*) *exhaustive*  
**begin**

**definition** *exhaustive-interval*::('a *interval* ⇒ (bool × term list) option)  
⇒ natural ⇒ (bool × term list) option

**where**  
*exhaustive-interval* f d =  
Quickcheck-Exhaustive.exhaustive (λx. Quickcheck-Exhaustive.exhaustive (λy.  
f (Ivl x y)) d) d

**instance** ⟨*proof*⟩

**end**

**definition** (in *term-syntax*) [*code-unfold*]:  
*valtermify-interval* x y = Code-Evaluation.valtermify (Ivl::'a::{preorder,typerep}⇒-)  
{·} x {·} y

**instantiation** *interval* :: (*{full-exhaustive,preorder,typerep}*) *full-exhaustive*  
**begin**

**definition** *full-exhaustive-interval*::  
(*'a interval* × (unit ⇒ term) ⇒ (bool × term list) option)  
⇒ natural ⇒ (bool × term list) option **where**  
*full-exhaustive-interval* f d =  
Quickcheck-Exhaustive.full-exhaustive  
(λx. Quickcheck-Exhaustive.full-exhaustive (λy. f (valtermify-interval x y)) d)  
d

**instance** ⟨*proof*⟩

**end**

**instantiation** *interval* :: (*{random,preorder,typerep}*) *random*  
**begin**

**definition** *random-interval* ::  
natural  
⇒ natural × natural  
⇒ ('a *interval* × (unit ⇒ term)) × natural × natural **where**  
*random-interval* i =  
scomp (Quickcheck-Random.random i)  
(λman. scomp (Quickcheck-Random.random i) (λexp. Pair (valtermify-interval  
man exp)))

**instance** ⟨*proof*⟩

**end**

end

### 3 Approximate Operations on Intervals of Floating Point Numbers

```
theory Interval-Approximation
  imports
    HOL-Decision-Proc.Approximation-Bounds
    Interval
begin
```

```
lifting-update float.lifting — TODO: in Float!
lifting-forget float.lifting
```

TODO: many of the lemmas should move to theories Float or Approximation (the latter should be based on type *interval*).

#### 3.1 Intervals with Floating Point Bounds

```
lift-definition round-interval :: nat  $\Rightarrow$  float interval  $\Rightarrow$  float interval
  is  $\lambda p. \lambda(l, u). (float-round-down\ p\ l, float-round-up\ p\ u)$ 
  <proof>
```

```
lemma lower-round-ivl[simp]: lower (round-interval p x) = float-round-down p
(lower x)
  <proof>
```

```
lemma upper-round-ivl[simp]: upper (round-interval p x) = float-round-up p (upper
x)
  <proof>
```

```
lemma round-ivl-correct: set-of A  $\subseteq$  set-of (round-interval prec A)
  <proof>
```

```
lift-definition truncate-ivl :: nat  $\Rightarrow$  real interval  $\Rightarrow$  real interval
  is  $\lambda p. \lambda(l, u). (truncate-down\ p\ l, truncate-up\ p\ u)$ 
  <proof>
```

```
lemma lower-truncate-ivl[simp]: lower (truncate-ivl p x) = truncate-down p (lower
x)
  <proof>
```

```
lemma upper-truncate-ivl[simp]: upper (truncate-ivl p x) = truncate-up p (upper
x)
  <proof>
```

```
lemma truncate-ivl-correct: set-of A  $\subseteq$  set-of (truncate-ivl prec A)
  <proof>
```

**lift-definition** *real-interval::float interval*  $\Rightarrow$  *real interval*  
**is**  $\lambda(l, u).$  (*real-of-float l, real-of-float u*)  
 $\langle$ *proof* $\rangle$

**lemma** *lower-real-interval[simp]*: *lower (real-interval x) = lower x*  
 $\langle$ *proof* $\rangle$

**lemma** *upper-real-interval[simp]*: *upper (real-interval x) = upper x*  
 $\langle$ *proof* $\rangle$

**definition** *set-of' x = (case x of None  $\Rightarrow$  UNIV | Some i  $\Rightarrow$  set-of (real-interval i))*

**lemma** *real-interval-min-interval[simp]*:  
*real-interval (min-interval a b) = min-interval (real-interval a) (real-interval b)*  
 $\langle$ *proof* $\rangle$

**lemma** *real-interval-max-interval[simp]*:  
*real-interval (max-interval a b) = max-interval (real-interval a) (real-interval b)*  
 $\langle$ *proof* $\rangle$

### 3.2 Intervals for standard functions

**lift-definition** *power-float-interval :: nat  $\Rightarrow$  nat  $\Rightarrow$  float interval  $\Rightarrow$  float interval*  
**is**  $\lambda p n (l, u).$  *float-power-bnds p n l u*  
 $\langle$ *proof* $\rangle$

**lemma** *lower-power-float-interval[simp]*:  
*lower (power-float-interval p n x) = fst (float-power-bnds p n (lower x) (upper x))*  
 $\langle$ *proof* $\rangle$

**lemma** *upper-power-float-interval[simp]*:  
*upper (power-float-interval p n x) = snd (float-power-bnds p n (lower x) (upper x))*  
 $\langle$ *proof* $\rangle$

**lemma** *power-float-intervalI*:  $x \in_i$  *real-interval X*  $\Longrightarrow$   $x \wedge n \in_i$  *real-interval (power-float-interval p n X)*  
 $\langle$ *proof* $\rangle$

**lift-definition** *mult-float-interval::nat  $\Rightarrow$  float interval  $\Rightarrow$  float interval  $\Rightarrow$  float interval*  
**is**  $\lambda prec.$   $\lambda(a1, a2).$   $\lambda(b1, b2).$  *bnds-mult prec a1 a2 b1 b2*  
 $\langle$ *proof* $\rangle$

**lemma** *lower-mult-float-interval[simp]*:  
*lower (mult-float-interval p x y) = fst (bnds-mult p (lower x) (upper x) (lower y) (upper y))*  
 $\langle$ *proof* $\rangle$

**lemma** *upper-mult-float-interval[simp]*:

$upper (mult\text{-}float\text{-}interval\ p\ x\ y) = snd (bnds\text{-}mult\ p\ (lower\ x)\ (upper\ x)\ (lower\ y)\ (upper\ y))$   
 ⟨proof⟩

**lemma** *mult-float-interval*:  
 $set\text{-}of\ (real\text{-}interval\ A) * set\text{-}of\ (real\text{-}interval\ B) \subseteq$   
 $set\text{-}of\ (real\text{-}interval\ (mult\text{-}float\text{-}interval\ prec\ A\ B))$   
 ⟨proof⟩

**lemma** *mult-float-intervalI*:  
 $x * y \in_i (real\text{-}interval\ (mult\text{-}float\text{-}interval\ prec\ A\ B))$   
**if**  $x \in_i real\text{-}interval\ A\ y \in_i real\text{-}interval\ B$   
 ⟨proof⟩

**lift-definition** *sqrt-float-interval*:: $nat \Rightarrow float\ interval \Rightarrow float\ interval$   
**is**  $\lambda prec. \lambda (lx, ux). (lb\text{-}sqrt\ prec\ lx, ub\text{-}sqrt\ prec\ ux)$   
 ⟨proof⟩

**lemma** *lower-float-interval[simp]*:  $lower (sqrt\text{-}float\text{-}interval\ prec\ X) = lb\text{-}sqrt\ prec$   
 $(lower\ X)$   
 ⟨proof⟩

**lemma** *upper-float-interval[simp]*:  $upper (sqrt\text{-}float\text{-}interval\ prec\ X) = ub\text{-}sqrt\ prec$   
 $(upper\ X)$   
 ⟨proof⟩

**lemma** *sqrt-float-interval*:  
 $sqrt \text{ ' } set\text{-}of\ (real\text{-}interval\ X) \subseteq set\text{-}of\ (real\text{-}interval\ (sqrt\text{-}float\text{-}interval\ prec\ X))$   
 ⟨proof⟩

**lemma** *sqrt-float-intervalI*:  
 $sqrt\ x \in_i real\text{-}interval\ (sqrt\text{-}float\text{-}interval\ p\ X)$   
**if**  $x \in set\text{-}of\ (real\text{-}interval\ X)$   
 ⟨proof⟩

**lemmas**  $[simp\ del] = lb\text{-}arctan.simps\ ub\text{-}arctan.simps$

**lemma** *lb-arctan*:  $arctan (real\text{-}of\text{-}float\ x) \leq y \Longrightarrow real\text{-}of\text{-}float (lb\text{-}arctan\ prec\ x)$   
 $\leq y$   
**and** *ub-arctan*:  $y \leq arctan\ x \Longrightarrow y \leq ub\text{-}arctan\ prec\ x$   
**for**  $x::float$  **and**  $y::real$   
 ⟨proof⟩

**lift-definition** *arctan-float-interval* ::  $nat \Rightarrow float\ interval \Rightarrow float\ interval$   
**is**  $\lambda prec. \lambda (lx, ux). (lb\text{-}arctan\ prec\ lx, ub\text{-}arctan\ prec\ ux)$   
 ⟨proof⟩

**lemma** *lower-arctan-float-interval[simp]*:  $lower (arctan\text{-}float\text{-}interval\ p\ x) = lb\text{-}arctan$   
 $p (lower\ x)$

$\langle \text{proof} \rangle$   
**lemma** *upper-arctan-float-interval[simp]*:  $\text{upper} (\text{arctan-float-interval } p \ x) = \text{ub-arctan } p \ (\text{upper } x)$   
 $\langle \text{proof} \rangle$

**lemma** *arctan-float-interval*:  
 $\text{arctan} \text{ ' set-of } (\text{real-interval } x) \subseteq \text{set-of } (\text{real-interval } (\text{arctan-float-interval } p \ x))$   
 $\langle \text{proof} \rangle$

**lemma** *arctan-float-intervalI*:  
 $\text{arctan } x \in_i \text{real-interval } (\text{arctan-float-interval } p \ X)$   
**if**  $x \in \text{set-of } (\text{real-interval } X)$   
 $\langle \text{proof} \rangle$

**lemma** *bnds-cos-lower*:  $\bigwedge x. \text{real-of-float } xl \leq x \implies x \leq \text{real-of-float } xu \implies \cos x \leq y \implies \text{real-of-float } (\text{fst } (\text{bnds-cos } \text{prec } xl \ xu)) \leq y$   
**and** *bnds-cos-upper*:  $\bigwedge x. \text{real-of-float } xl \leq x \implies x \leq \text{real-of-float } xu \implies y \leq \cos x \implies y \leq \text{real-of-float } (\text{snd } (\text{bnds-cos } \text{prec } xl \ xu))$   
**for**  $xl \ xu :: \text{float}$  **and**  $y :: \text{real}$   
 $\langle \text{proof} \rangle$

**lift-definition** *cos-float-interval* ::  $\text{nat} \Rightarrow \text{float interval} \Rightarrow \text{float interval}$   
**is**  $\lambda \text{prec}. \lambda (lx, ux). \text{bnds-cos } \text{prec } lx \ ux$   
 $\langle \text{proof} \rangle$

**lemma** *lower-cos-float-interval[simp]*:  $\text{lower} (\text{cos-float-interval } p \ x) = \text{fst } (\text{bnds-cos } p \ (\text{lower } x) \ (\text{upper } x))$   
 $\langle \text{proof} \rangle$

**lemma** *upper-cos-float-interval[simp]*:  $\text{upper} (\text{cos-float-interval } p \ x) = \text{snd } (\text{bnds-cos } p \ (\text{lower } x) \ (\text{upper } x))$   
 $\langle \text{proof} \rangle$

**lemma** *cos-float-interval*:  
 $\text{cos} \text{ ' set-of } (\text{real-interval } x) \subseteq \text{set-of } (\text{real-interval } (\text{cos-float-interval } p \ x))$   
 $\langle \text{proof} \rangle$

**lemma** *cos-float-intervalI*:  
 $\text{cos } x \in_i \text{real-interval } (\text{cos-float-interval } p \ X)$   
**if**  $x \in \text{set-of } (\text{real-interval } X)$   
 $\langle \text{proof} \rangle$

**lemma** *lb-exp*:  $\text{exp } x \leq y \implies \text{lb-exp } \text{prec } x \leq y$   
**and** *ub-exp*:  $y \leq \text{exp } x \implies y \leq \text{ub-exp } \text{prec } x$   
**for**  $x :: \text{float}$  **and**  $y :: \text{real}$   $\langle \text{proof} \rangle$

**lift-definition** *exp-float-interval* ::  $\text{nat} \Rightarrow \text{float interval} \Rightarrow \text{float interval}$   
**is**  $\lambda \text{prec}. \lambda (lx, ux). (\text{lb-exp } \text{prec } lx, \text{ub-exp } \text{prec } ux)$   
 $\langle \text{proof} \rangle$

**lemma** *lower-exp-float-interval[simp]*:  $\text{lower } (\text{exp-float-interval } p \ x) = \text{lb-exp } p$   
*(lower x)*

*<proof>*

**lemma** *upper-exp-float-interval[simp]*:  $\text{upper } (\text{exp-float-interval } p \ x) = \text{ub-exp } p$   
*(upper x)*

*<proof>*

**lemma** *exp-float-interval*:

$\text{exp } \text{' set-of } (\text{real-interval } x) \subseteq \text{set-of } (\text{real-interval } (\text{exp-float-interval } p \ x))$

*<proof>*

**lemma** *exp-float-intervalI*:

$\text{exp } x \in_i \text{real-interval } (\text{exp-float-interval } p \ X)$

**if**  $x \in \text{set-of } (\text{real-interval } X)$

*<proof>*

**lemmas** *[simp del]* = *lb-ln.simps ub-ln.simps*

**lemma** *lb-lnD*:

$y \leq \text{ln } x \wedge 0 < \text{real-of-float } x$  **if**  $\text{lb-ln prec } x = \text{Some } y$

*<proof>*

**lemma** *ub-lnD*:

$\text{ln } x \leq y \wedge 0 < \text{real-of-float } x$  **if**  $\text{ub-ln prec } x = \text{Some } y$

*<proof>*

**lift-definition**(*code-dt*) *ln-float-interval* ::  $\text{nat} \Rightarrow \text{float interval} \Rightarrow \text{float interval option}$

**is**  $\lambda \text{prec. } \lambda (lx, ux).$

$\text{Option.bind } (\text{lb-ln prec } lx) (\lambda l.$

$\text{Option.bind } (\text{ub-ln prec } ux) (\lambda u. \text{Some } (l, u)))$

*<proof>*

**lemma** *ln-float-interval-eq-Some-conv[simp]*:

$\text{ln-float-interval } p \ x = \text{Some } y \iff$

$\text{lb-ln } p \ (\text{lower } x) = \text{Some } (\text{lower } y) \wedge \text{ub-ln } p \ (\text{upper } x) = \text{Some } (\text{upper } y)$

*<proof>*

**lemma** *ln-float-interval*:  $\text{ln } \text{' set-of } (\text{real-interval } x) \subseteq \text{set-of } (\text{real-interval } y)$

**if**  $\text{ln-float-interval } p \ x = \text{Some } y$

*<proof>*

**lemma** *ln-float-intervalI*:

$\text{ln } x \in \text{set-of}' (\text{ln-float-interval } p \ X)$

**if**  $x \in_i (\text{real-interval } X)$

*<proof>*

**lift-definition**(*code-dt*) *powr-float-interval* ::  $\text{nat} \Rightarrow \text{float interval} \Rightarrow \text{float interval} \Rightarrow \text{float interval option}$

**is**  $\lambda prec. \lambda(l1, u1). \lambda(l2, u2). bnds-powr\ prec\ l1\ u1\ l2\ u2$   
 ⟨proof⟩

**lemma** *powr-float-interval*:  
 $\{x\ powr\ y \mid x\ y. x \in set-of\ (real-interval\ X) \wedge y \in set-of\ (real-interval\ Y)\}$   
 $\subseteq set-of\ (real-interval\ R)$   
**if** *powr-float-interval*  $prec\ X\ Y = Some\ R$   
 ⟨proof⟩

**lemma** *powr-float-intervalI*:  
 $x\ powr\ y \in set-of'\ (powr-float-interval\ p\ X\ Y)$   
**if**  $x \in_i\ real-interval\ X\ y \in_i\ real-interval\ Y$   
 ⟨proof⟩

**lift-definition**(*code-dt*) *inverse-float-interval::nat*  $\Rightarrow$  *float interval*  $\Rightarrow$  *float interval*  
**option is**  
 $\lambda prec\ (l, u). if\ (0 < l \vee u < 0)\ then\ Some\ (float-divl\ prec\ 1\ u, float-divr\ prec\ 1\ l)$   
 else *None*  
 ⟨proof⟩

**lemma** *inverse-float-interval-eq-Some-conv[simp]*:  
**defines**  $one \equiv (1::float)$   
**shows**  
 $inverse-float-interval\ p\ X = Some\ R \iff$   
 $(lower\ X > 0 \vee upper\ X < 0) \wedge$   
 $lower\ R = float-divl\ p\ one\ (upper\ X) \wedge$   
 $upper\ R = float-divr\ p\ one\ (lower\ X)$   
 ⟨proof⟩

**lemma** *inverse-float-interval*:  
 $inverse\ 'set-of\ (real-interval\ X) \subseteq set-of\ (real-interval\ Y)$   
**if** *inverse-float-interval*  $p\ X = Some\ Y$   
 ⟨proof⟩

**lemma** *inverse-float-intervalI*:  
 $x \in set-of\ (real-interval\ X) \implies inverse\ x \in set-of'\ (inverse-float-interval\ p\ X)$   
 ⟨proof⟩

**lift-definition** *pi-float-interval::nat*  $\Rightarrow$  *float interval* **is**  $\lambda prec. (lb-pi\ prec, ub-pi\ prec)$   
 ⟨proof⟩

**lemma** *lower-pi-float-interval[simp]*:  $lower\ (pi-float-interval\ prec) = lb-pi\ prec$   
 ⟨proof⟩

**lemma** *upper-pi-float-interval[simp]*:  $upper\ (pi-float-interval\ prec) = ub-pi\ prec$   
 ⟨proof⟩

**lemma** *pi-float-interval*:  $pi \in set-of\ (real-interval\ (pi-float-interval\ prec))$   
 ⟨proof⟩

**lemma** *real-interval-abs-interval*[simp]:  
*real-interval* (abs-interval  $x$ ) = abs-interval (real-interval  $x$ )  
 ⟨proof⟩

**lift-definition** *floor-float-interval*::float interval  $\Rightarrow$  float interval **is**  
 $\lambda(l, u).$  (floor-fl  $l$ , floor-fl  $u$ )  
 ⟨proof⟩

**lemma** *lower-floor-float-interval*[simp]: lower (floor-float-interval  $x$ ) = floor-fl (lower  $x$ )  
 ⟨proof⟩

**lemma** *upper-floor-float-interval*[simp]: upper (floor-float-interval  $x$ ) = floor-fl (upper  $x$ )  
 ⟨proof⟩

**lemma** *floor-float-intervalI*:  $[x] \in_i$  real-interval (floor-float-interval  $X$ )  
**if**  $x \in_i$  real-interval  $X$   
 ⟨proof⟩

**lemma** *in-intervalI*:  
 $x \in_i X$  **if** lower  $X \leq x \leq$  upper  $X$   
 ⟨proof⟩

**abbreviation** *in-real-interval* ((-/  $\in_r$  -) [51, 51] 50) **where**  
 $x \in_r X \equiv x \in_i$  real-interval  $X$

**lemma** *in-real-intervalI*:  
 $x \in_r X$  **if** lower  $X \leq x \leq$  upper  $X$  **for**  $x::$ real **and**  $X::$ float interval  
 ⟨proof⟩

**lemma** *lower-Interval*: lower (Interval  $x$ ) = fst  $x$   
**and** *upper-Interval*: upper (Interval  $x$ ) = snd  $x$   
**if** fst  $x \leq$  snd  $x$   
 ⟨proof⟩

**definition** *all-in-i* :: 'a::preorder list  $\Rightarrow$  'a interval list  $\Rightarrow$  bool  
 (infix (all'-in<sub>*i*</sub>) 50)  
**where**  $x$  all-in<sub>*i*</sub>  $I =$  (length  $x =$  length  $I \wedge (\forall i <$  length  $I. x ! i \in_i I ! i)$ )

**definition** *all-in* :: real list  $\Rightarrow$  float interval list  $\Rightarrow$  bool  
 (infix (all'-in) 50)  
**where**  $x$  all-in  $I =$  (length  $x =$  length  $I \wedge (\forall i <$  length  $I. x ! i \in_r I ! i)$ )

**definition** *all-subset* :: 'a::order interval list  $\Rightarrow$  'a interval list  $\Rightarrow$  bool  
 (infix (all'-subset) 50)  
**where**  $I$  all-subset  $J =$  (length  $I =$  length  $J \wedge (\forall i <$  length  $I. \text{set-of } (I!i) \subseteq \text{set-of } (J!i))$ )

**lemmas** [simp] = all-in-def all-subset-def

**lemma** all-subsetD:

**assumes**  $I$  all-subset  $J$

**assumes**  $x$  all-in  $I$

**shows**  $x$  all-in  $J$

*<proof>*

**lemma** plus-down-mono: plus-down  $p$   $a$   $b$   $\leq$  plus-down  $p$   $c$   $d$  **if**  $a + b \leq c + d$

*<proof>*

**lemma** plus-up-mono: plus-up  $p$   $a$   $b$   $\leq$  plus-up  $p$   $c$   $d$  **if**  $a + b \leq c + d$

*<proof>*

**lemma** round-interval-mono: set-of (round-interval prec  $X$ )  $\subseteq$  set-of (round-interval prec  $Y$ )

**if** set-of  $X$   $\subseteq$  set-of  $Y$

*<proof>*

**lemma** mult-mono-nonpos-nonneg:  $a * b \leq c * d$

**if**  $a \leq c$   $a \leq 0$   $0 \leq d$   $d \leq b$  **for**  $a$   $b$   $c$   $d$ ::'a::ordered-ring

*<proof>*

**lemma** mult-mono-nonneg-nonpos:  $b * a \leq d * c$

**if**  $a \leq c$   $c \leq 0$   $0 \leq d$   $d \leq b$  **for**  $a$   $b$   $c$   $d$ ::'a::ordered-ring

*<proof>*

**lemma** mult-mono-nonpos-nonpos:  $a * b \leq c * d$

**if**  $a \geq c$   $a \leq 0$   $b \geq d$   $d \leq 0$  **for**  $a$   $b$   $c$   $d$ ::real

*<proof>*

**lemma** mult-float-mono1:

**notes** mono-rules = plus-down-mono add-mono nprt-mono nprt-le-zero zero-le-pprt pprrt-mono

**shows**  $a \leq b \implies ab \leq bb \implies$

$aa \leq a \implies$

$b \leq ba \implies$

$ac \leq ab \implies$

$bb \leq bc \implies$

plus-down prec (nprt  $aa$  \* pprrt  $bc$ )

(plus-down prec (nprt  $ba$  \* nprt  $bc$ )

(plus-down prec (pprrt  $aa$  \* pprrt  $ac$ )

(pprrt  $ba$  \* nprt  $ac$ )))

$\leq$  plus-down prec (nprt  $a$  \* pprrt  $bb$ )

(plus-down prec (nprt  $b$  \* nprt  $bb$ )

(plus-down prec (pprrt  $a$  \* pprrt  $ab$ )

(pprrt  $b$  \* nprt  $ab$ )))

*<proof>*

**lemma** *mult-float-mono2*:

**notes** *mono-rules* = *plus-up-mono add-mono nprt-mono nprt-le-zero zero-le-pprt pprt-mono*

**shows**  $a \leq b \implies$

$ab \leq bb \implies$

$aa \leq a \implies$

$b \leq ba \implies$

$ac \leq ab \implies$

$bb \leq bc \implies$

*plus-up prec (pprt b \* pprt bb)*

*(plus-up prec (pprt a \* nprt bb)*

*(plus-up prec (nprt b \* pprt ab)*

*(nprt a \* nprt ab)))*

$\leq$  *plus-up prec (pprt ba \* pprt bc)*

*(plus-up prec (pprt aa \* nprt bc)*

*(plus-up prec (nprt ba \* pprt ac)*

*(nprt aa \* nprt ac)))*

*<proof>*

**lemma** *mult-float-interval-mono*: *set-of (mult-float-interval prec A B)  $\subseteq$  set-of (mult-float-interval prec X Y)*

**if** *set-of A  $\subseteq$  set-of X set-of B  $\subseteq$  set-of Y*

*<proof>*

**lemma** *Ivl-simps[simp]*: *lower (Ivl a b) = min a b upper (Ivl a b) = b*

*<proof>*

**lemmas** [*simp del*] = *power-down.simps(2) power-up.simps(2)*

**lemmas** *power-down-simp = power-down.simps(2)*

**lemmas** *power-up-simp = power-up.simps(2)*

**lemma** *power-down-even-nonneg*: *even n  $\implies 0 \leq$  power-down p x n*

*<proof>*

**lemma** *truncate-down-less-zero-iff[simp]*: *truncate-down p x < 0  $\iff$  x < 0*

*<proof>*

**lemma** *truncate-down-nonneg-iff[simp]*: *truncate-down p x  $\geq 0 \iff$  x  $\geq 0$*

*<proof>*

**lemma** *truncate-down-eq-zero-iff[simp]*: *truncate-down prec x = 0  $\iff$  x = 0*

*<proof>*

**lemma** *power-down-eq-zero-iff[simp]*: *power-down prec b n = 0  $\iff$  b = 0  $\wedge$  n  $\neq$  0*

*<proof>*

**lemma** *power-down-nonneg-iff[simp]*:

*power-down prec b n*  $\geq 0 \iff \text{even } n \vee b \geq 0$   
(proof)

**lemma** *power-down-neg-iff*[simp]:

*power-down prec b n*  $< 0 \iff$   
 $b < 0 \wedge \text{odd } n$

(proof)

**lemma** *power-down-nonpos-iff*[simp]:

**notes** [simp del] = *power-down-neg-iff power-down-eq-zero-iff*

**shows** *power-down prec b n*  $\leq 0 \iff b < 0 \wedge \text{odd } n \vee b = 0 \wedge n \neq 0$

(proof)

**lemma** *power-down-mono*:

*power-down prec a n*  $\leq$  *power-down prec b n*

**if**  $((0 \leq a \wedge a \leq b) \vee (\text{odd } n \wedge a \leq b) \vee (\text{even } n \wedge a \leq 0 \wedge b \leq a))$

(proof)

**lemma** *truncate-up-nonneg*:  $0 \leq \text{truncate-up } p \ x$  **if**  $0 \leq x$

(proof)

**lemma** *truncate-up-pos*:  $0 < \text{truncate-up } p \ x$  **if**  $0 < x$

(proof)

**lemma** *truncate-up-less-zero-iff*[simp]: *truncate-up p x*  $< 0 \iff x < 0$

(proof)

**lemma** *truncate-up-nonneg-iff*[simp]: *truncate-up p x*  $\geq 0 \iff x \geq 0$

(proof)

**lemma** *power-up-even-nonneg*:  $\text{even } n \implies 0 \leq \text{power-up } p \ x \ n$

(proof)

**lemma** *truncate-up-eq-zero-iff*[simp]: *truncate-up prec x*  $= 0 \iff x = 0$

(proof)

**lemma** *power-up-eq-zero-iff*[simp]: *power-up prec b n*  $= 0 \iff b = 0 \wedge n \neq 0$

(proof)

**lemma** *power-up-nonneg-iff*[simp]:

*power-up prec b n*  $\geq 0 \iff \text{even } n \vee b \geq 0$

(proof)

**lemma** *power-up-neg-iff*[simp]:

*power-up prec b n*  $< 0 \iff b < 0 \wedge \text{odd } n$

(proof)

**lemma** *power-up-nonpos-iff*[simp]:

**notes** [simp del] = *power-up-neg-iff power-up-eq-zero-iff*

**shows**  $\text{power-up prec } b \ n \leq 0 \iff b < 0 \wedge \text{odd } n \vee b = 0 \wedge n \neq 0$   
 ⟨proof⟩

**lemma** *power-up-mono*:

$\text{power-up prec } a \ n \leq \text{power-up prec } b \ n$   
**if**  $((0 \leq a \wedge a \leq b) \vee (\text{odd } n \wedge a \leq b) \vee (\text{even } n \wedge a \leq 0 \wedge b \leq a))$   
 ⟨proof⟩

**lemma** *set-of-subset-iff*:  $\text{set-of } X \subseteq \text{set-of } Y \iff \text{lower } Y \leq \text{lower } X \wedge \text{upper } X$   
 $\leq \text{upper } Y$

**for**  $X \ Y :: 'a :: \text{linorder interval}$   
 ⟨proof⟩

**lemma** *power-float-interval-mono*:

$\text{set-of } (\text{power-float-interval prec } n \ A)$   
 $\subseteq \text{set-of } (\text{power-float-interval prec } n \ B)$   
**if**  $\text{set-of } A \subseteq \text{set-of } B$   
 ⟨proof⟩

**lemma** *bounds-of-interval-eq-lower-upper*:

$\text{bounds-of-interval } \text{ivl} = (\text{lower } \text{ivl}, \text{upper } \text{ivl})$  **if**  $\text{lower } \text{ivl} \leq \text{upper } \text{ivl}$   
 ⟨proof⟩

**lemma** *real-interval-Ivl*:  $\text{real-interval } (\text{Ivl } a \ b) = \text{Ivl } a \ b$

⟨proof⟩

**lemma** *set-of-mul-contains-real-zero*:

$0 \in_r (A * B)$  **if**  $0 \in_r A \vee 0 \in_r B$   
 ⟨proof⟩

**fun** *subdivide-interval* ::  $\text{nat} \Rightarrow \text{float interval} \Rightarrow \text{float interval list}$

**where** *subdivide-interval*  $0 \ I = [I]$   
 | *subdivide-interval*  $(\text{Suc } n) \ I =$   
    $\text{let } m = \text{mid } I$   
    $\text{in } (\text{subdivide-interval } n \ (\text{Ivl } (\text{lower } I) \ m)) @ (\text{subdivide-interval } n \ (\text{Ivl } m$   
 $(\text{upper } I)))$   
 )

**lemma** *subdivide-interval-length*:

**shows**  $\text{length } (\text{subdivide-interval } n \ I) = 2^n$   
 ⟨proof⟩

**lemma** *lower-le-mid*:  $\text{lower } x \leq \text{mid } x$  *real-of-float*  $(\text{lower } x) \leq \text{mid } x$

**and** *mid-le-upper*:  $\text{mid } x \leq \text{upper } x$  *real-of-float*  $(\text{mid } x) \leq \text{upper } x$   
 ⟨proof⟩

**lemma** *subdivide-interval-correct*:

*list-ex*  $(\lambda i. x \in_r i) (\text{subdivide-interval } n \ I)$  **if**  $x \in_r I$  **for**  $x :: \text{real}$   
 ⟨proof⟩

```

fun interval-list-union :: 'a::lattice interval list  $\Rightarrow$  'a interval
  where interval-list-union [] = undefined
  | interval-list-union [I] = I
  | interval-list-union (I#Is) = sup I (interval-list-union Is)

```

```

lemma interval-list-union-correct:
  assumes S  $\neq$  []
  assumes i < length S
  shows set-of (S!i)  $\subseteq$  set-of (interval-list-union S)
  <proof>

```

```

lemma split-domain-correct:
  fixes x :: real list
  assumes x all-in I
  assumes split-correct:  $\bigwedge x a I. x \in_r I \implies$  list-ex ( $\lambda i::\text{float interval. } x \in_r i$ ) (split I)
  shows list-ex ( $\lambda s. x$  all-in s) (split-domain split I)
  <proof>

```

**end**

## 4 Horner Evaluation

```

theory Horner-Eval
  imports Interval
begin

```

Function and lemmas for evaluating polynomials via the horner scheme. Because interval multiplication is not distributive, interval polynomials expressed as a sum of monomials are not equivalent to their respective horner form. The functions and lemmas in this theory can be used to express interval polynomials in horner form and prove facts about them.

```

fun horner-eval' where
  horner-eval' f x v 0 = v
  | horner-eval' f x v (Suc i) = horner-eval' f x (f i + x * v) i

```

```

definition horner-eval
  where horner-eval f x n = horner-eval' f x 0 n

```

```

lemma horner-eval-cong:
  assumes  $\bigwedge i. i < n \implies f i = g i$ 
  assumes x = y
  assumes n = m
  shows horner-eval f x n = horner-eval g y m
  <proof>

```

```

lemma horner-eval-eq-setsum:

```

```

fixes  $x::'a::\text{linordered-idom}$ 
shows  $\text{horner-eval } f \ x \ n = (\sum i < n. f \ i \ * \ x^i)$ 
<proof>

lemma  $\text{horner-eval-Suc}[simp]$ :
fixes  $x::'a::\text{linordered-idom}$ 
shows  $\text{horner-eval } f \ x \ (\text{Suc } n) = \text{horner-eval } f \ x \ n + (f \ n) * x^n$ 
<proof>

lemma  $\text{horner-eval-Suc}'[simp]$ :
fixes  $x::'a::\{\text{comm-monoid-add}, \text{times}\}$ 
shows  $\text{horner-eval } f \ x \ (\text{Suc } n) = f \ 0 + x * (\text{horner-eval } (\lambda i. f \ (\text{Suc } i)) \ x \ n)$ 
<proof>

lemma  $\text{horner-eval-0}[simp]$ :
shows  $\text{horner-eval } f \ x \ 0 = 0$ 
<proof>

lemma  $\text{horner-eval}'\text{-interval}$ :
fixes  $x::'a::\text{linordered-ring}$ 
assumes  $\bigwedge i. i < n \implies f \ i \in \text{set-of } (g \ i)$ 
assumes  $x \in_i I \ v \in_i V$ 
shows  $\text{horner-eval}' \ f \ x \ v \ n \in_i \text{horner-eval}' \ g \ I \ V \ n$ 
<proof>

lemma  $\text{horner-eval-interval}$ :
fixes  $x::'a::\text{linordered-idom}$ 
assumes  $\bigwedge i. i < n \implies f \ i \in \text{set-of } (g \ i)$ 
assumes  $x \in \text{set-of } I$ 
shows  $\text{horner-eval } f \ x \ n \in_i \text{horner-eval } g \ I \ n$ 
<proof>

end
theory  $\text{Polynomial-Expression-Additional}$ 
imports  $\text{Interval-Approximation}$ 
 $\text{Polynomial-Expression}$ 
 $\text{HOL-Decision-Procs.Approximation}$ 
begin

lemma  $\text{real-of-float-eq-zero-iff}[simp]$ :  $\text{real-of-float } x = 0 \iff x = 0$ 
<proof>

Theory  $\text{Taylor-Models.Polynomial-Expression}$  contains a, more or less, 1:1
generalization of theory  $\text{Multivariate-Polynomial}$ . Any additions belong
here.

declare  $[[\text{coercion-map } \text{map-poly}]]$ 
declare  $[[\text{coercion } \text{interval-of}::\text{float} \Rightarrow \text{float } \text{interval}]]$ 

Apply float interval arguments to a float poly.

```

**value** *Ipoly* [*Ivl* (*Float* 4 (-6)) (*Float* 10 6)] (*poly.Add* (*poly.C* (*Float* 3 5))  
(*poly.Bound* 0))

*map-poly* for homomorphisms

**lemma** *map-poly-homo-polyadd-eq-zero-iff*:

$$\text{map-poly } f (p +_p q) = 0_p \longleftrightarrow p +_p q = 0_p$$

$$\text{if } [\text{symmetric}, \text{simp}]: \bigwedge x y. f (x + y) = f x + f y \bigwedge x. f x = 0 \longleftrightarrow x = 0$$

*<proof>*

**lemma** *zero-iffD*:  $(\bigwedge x. f x = 0 \longleftrightarrow x = 0) \implies f 0 = 0$

*<proof>*

**lemma** *map-poly-homo-polyadd*:

$$\text{map-poly } f (p1 +_p p2) = \text{map-poly } f p1 +_p \text{map-poly } f p2$$

$$\text{if } [\text{simp}]: \bigwedge x y. f (x + y) = f x + f y \bigwedge x. f x = 0 \longleftrightarrow x = 0$$

*<proof>*

**lemma** *map-poly-homo-polyneg*:

$$\text{map-poly } f (\sim_p p1) = \sim_p (\text{map-poly } f p1)$$

$$\text{if } [\text{simp}]: \bigwedge x y. f (-x) = -f x$$

*<proof>*

**lemma** *map-poly-homo-polysub*:

$$\text{map-poly } f (p1 -_p p2) = \text{map-poly } f p1 -_p \text{map-poly } f p2$$

$$\text{if } [\text{simp}]: \bigwedge x y. f (x + y) = f x + f y \bigwedge x. f x = 0 \longleftrightarrow x = 0 \bigwedge x y. f (-x) = -f x$$

*<proof>*

**lemma** *map-poly-homo-polymul*:

$$\text{map-poly } f (p1 *_p p2) = \text{map-poly } f p1 *_p \text{map-poly } f p2$$

$$\text{if } [\text{simp}]: \bigwedge x y. f (x + y) = f x + f y \bigwedge x. f x = 0 \longleftrightarrow x = 0 \bigwedge x y. f (x * y) = f x * f y$$

*<proof>*

**lemma** *map-poly-homo-polypow*:

$$\text{map-poly } f (p1 \hat{^}_p n) = \text{map-poly } f p1 \hat{^}_p n$$

$$\text{if } [\text{simp}]: \bigwedge x y. f (x + y) = f x + f y \bigwedge x. f x = 0 \longleftrightarrow x = 0 \bigwedge x y. f (x * y) = f x * f y$$

$$f 1 = 1$$

*<proof>*

**lemmas** *map-poly-homo-polyarith* = *map-poly-homo-polyadd* *map-poly-homo-polyneg*  
*map-poly-homo-polysub* *map-poly-homo-polymul* *map-poly-homo-polypow*

Count the number of parameters of a polynomial.

**fun** *num-params* :: 'a *poly*  $\Rightarrow$  *nat*

**where** *num-params* (*poly.C* *c*) = 0

| *num-params* (*poly.Bound* *n*) = *Suc* *n*

| *num-params* (*poly.Add* *a* *b*) = *max* (*num-params* *a*) (*num-params* *b*)

```

| num-params (poly.Sub a b) = max (num-params a) (num-params b)
| num-params (poly.Mul a b) = max (num-params a) (num-params b)
| num-params (poly.Neg a)   = num-params a
| num-params (poly.Pw a n)  = num-params a
| num-params (poly.CN a n b) = max (max (num-params a) (num-params b))
(Suc n)

```

**lemma** *num-params-map-poly[simp]*:  
**shows**  $\text{num-params } (\text{map-poly } f \ p) = \text{num-params } p$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-polyadd*:  
**shows**  $\text{num-params } (p1 \ +_p \ p2) \leq \max (\text{num-params } p1) (\text{num-params } p2)$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-polyneg*:  
**shows**  $\text{num-params } (\sim_p \ p) = \text{num-params } p$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-polymul*:  
**shows**  $\text{num-params } (p1 \ *_p \ p2) \leq \max (\text{num-params } p1) (\text{num-params } p2)$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-polypow*:  
**shows**  $\text{num-params } (p \ \hat{^}_p \ n) \leq \text{num-params } p$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-polynate*:  
**shows**  $\text{num-params } (\text{polynate } p) \leq \text{num-params } p$   
 $\langle \text{proof} \rangle$

**lemma** *polynate-map-poly-real[simp]*:  
**fixes**  $p :: \text{float poly}$   
**shows**  $\text{map-poly real-of-float } (\text{polynate } p) = \text{polynate } (\text{map-poly real-of-float } p)$   
 $\langle \text{proof} \rangle$

Evaluating a float poly is equivalent to evaluating the corresponding real poly with the float parameters converted to reals.

**lemma** *Ipoly-real-float-equiv*:  
**fixes**  $p :: \text{float poly}$  **and**  $xs :: \text{float list}$   
**assumes**  $\text{num-params } p \leq \text{length } xs$   
**shows**  $\text{Ipoly } xs \ (p :: \text{real poly}) = \text{Ipoly } xs \ p$   
 $\langle \text{proof} \rangle$

Evaluating an 'a poly with 'a interval arguments is monotone.

**lemma** *Ipoly-interval-args-mono*:  
**fixes**  $p :: 'a :: \text{linordered-idom poly}$   
**and**  $x :: 'a \text{ list}$   
**and**  $xs :: 'a \text{ interval list}$

**assumes**  $x$  all-in <sub>$i$</sub>   $xs$   
**assumes** num-params  $p \leq \text{length } xs$   
**shows**  $\text{Ipoly } x \ p \in \text{set-of } (\text{Ipoly } xs \ (\text{map-poly interval-of } p))$   
 $\langle \text{proof} \rangle$

**lemma** *Ipoly-interval-args-inc-mono*:  
**fixes**  $p::'a::\{\text{real-normed-algebra}, \text{linear-continuum-topology}, \text{linordered-idom}\}$   $\text{poly}$   
**and**  $I::'a$  interval list **and**  $J::'a$  interval list  
**assumes** num-params  $p \leq \text{length } I$   
**assumes**  $I$  all-subset  $J$   
**shows**  $\text{set-of } (\text{Ipoly } I \ (\text{map-poly interval-of } p)) \subseteq \text{set-of } (\text{Ipoly } J \ (\text{map-poly interval-of } p))$   
 $\langle \text{proof} \rangle$

## 5 Splitting polynomials to reduce floating point precision

TODO: Move this! Definitions regarding floating point numbers should not be in a theory about polynomials.

**fun** *float-prec* ::  $\text{float} \Rightarrow \text{int}$   
**where** *float-prec*  $f = (\text{let } p = \text{exponent } f \text{ in if } p \geq 0 \text{ then } 0 \text{ else } -p)$

**fun** *float-round* ::  $\text{nat} \Rightarrow \text{float} \Rightarrow \text{float}$   
**where** *float-round*  $\text{prec } f =$   
 $\text{let } d = \text{float-down } \text{prec } f; u = \text{float-up } \text{prec } f$   
 $\text{in if } f - d < u - f \text{ then } d \text{ else } u$

Splits any polynomial  $p$  into two polynomials  $l, r$ , such that  $\forall x::\text{real}. p(x) = l(x) + r(x)$  and all floating point coefficients in  $p$  are rounded to precision  $\text{prec}$ . Not all cases need to give good results. Polynomials normalized with `polynat` only contain `poly.C` and `poly.CN` constructors.

**fun** *split-by-prec* ::  $\text{nat} \Rightarrow \text{float } \text{poly} \Rightarrow \text{float } \text{poly} * \text{float } \text{poly}$   
**where** *split-by-prec*  $\text{prec } (\text{poly.C } f) = (\text{let } r = \text{float-round } \text{prec } f \text{ in } (\text{poly.C } r, \text{poly.C } (f - r)))$   
 $| \text{split-by-prec } \text{prec } (\text{poly.Bound } n) = (\text{poly.Bound } n, \text{poly.C } 0)$   
 $| \text{split-by-prec } \text{prec } (\text{poly.Add } l \ r) = (\text{let } (ll, lr) = \text{split-by-prec } \text{prec } l;$   
 $\quad (rl, rr) = \text{split-by-prec } \text{prec } r$   
 $\quad \text{in } (\text{poly.Add } ll \ rl, \text{poly.Add } lr \ rr))$   
 $| \text{split-by-prec } \text{prec } (\text{poly.Sub } l \ r) = (\text{let } (ll, lr) = \text{split-by-prec } \text{prec } l;$   
 $\quad (rl, rr) = \text{split-by-prec } \text{prec } r$   
 $\quad \text{in } (\text{poly.Sub } ll \ rl, \text{poly.Sub } lr \ rr))$   
 $| \text{split-by-prec } \text{prec } (\text{poly.Mul } l \ r) = (\text{let } (ll, lr) = \text{split-by-prec } \text{prec } l;$   
 $\quad (rl, rr) = \text{split-by-prec } \text{prec } r$   
 $\quad \text{in } (\text{poly.Mul } ll \ rl, \text{poly.Add } (\text{poly.Add } (\text{poly.Mul } lr \ rl) (\text{poly.Mul } ll \ rr)) (\text{poly.Mul } lr \ rr)))$   
 $| \text{split-by-prec } \text{prec } (\text{poly.Neg } p) = (\text{let } (l, r) = \text{split-by-prec } \text{prec } p \text{ in } (\text{poly.Neg } l, \text{poly.Neg } r))$

```

| split-by-prec prec (poly.Pw p 0) = (poly.C 1, poly.C 0)
| split-by-prec prec (poly.Pw p (Suc n)) = (let (l, r) = split-by-prec prec p in
(poly.Pw l n, poly.Sub (poly.Pw p (Suc n)) (poly.Pw l n)))
| split-by-prec prec (poly.CN c n p) = (let (cl, cr) = split-by-prec prec c;
(pl, pr) = split-by-prec prec p
in (poly.CN cl n pl, poly.CN cr n pr))

```

TODO: Prove precision constraint on  $l$ .

**lemma** *split-by-prec-correct*:

```

fixes args :: real list
assumes (l, r) = split-by-prec prec p
shows Ipoly args p = Ipoly args l + Ipoly args r (is ?P1)
and num-params l ≤ num-params p (is ?P2)
and num-params r ≤ num-params p (is ?P3)
⟨proof⟩

```

## 6 Splitting polynomials by degree

**fun** *maxdegree* :: ('a::zero) poly ⇒ nat

```

where maxdegree (poly.C c) = 0
| maxdegree (poly.Bound n) = 1
| maxdegree (poly.Add l r) = max (maxdegree l) (maxdegree r)
| maxdegree (poly.Sub l r) = max (maxdegree l) (maxdegree r)
| maxdegree (poly.Mul l r) = maxdegree l + maxdegree r
| maxdegree (poly.Neg p) = maxdegree p
| maxdegree (poly.Pw p n) = n * maxdegree p
| maxdegree (poly.CN c n p) = max (maxdegree c) (1 + maxdegree p)

```

**fun** *split-by-degree* :: nat ⇒ 'a::zero poly ⇒ 'a poly \* 'a poly

```

where split-by-degree n (poly.C c) = (poly.C c, poly.C 0)
| split-by-degree 0 p = (poly.C 0, p)
| split-by-degree (Suc n) (poly.CN c v p) = (
let (cl, cr) = split-by-degree (Suc n) c;
(pl, pr) = split-by-degree n p
in (poly.CN cl v pl, poly.CN cr v pr))

```

— This function is only intended for use on polynomials in normal form. Hence most cases never get executed.

```

| split-by-degree n p = (poly.C 0, p)

```

**lemma** *split-by-degree-correct*:

```

fixes x :: real list and p :: float poly
assumes (l, r) = split-by-degree ord p
shows maxdegree l ≤ ord (is ?P1)
and Ipoly x p = Ipoly x l + Ipoly x r (is ?P2)
and num-params l ≤ num-params p (is ?P3)
and num-params r ≤ num-params p (is ?P4)
⟨proof⟩

```

Operations on lists.

**lemma** *length-map2[simp]*:  $\text{length } (\text{map2 } f \ a \ b) = \min (\text{length } a) (\text{length } b)$   
 ⟨*proof*⟩

**lemma** *map2-nth[simp]*:  
**assumes**  $n < \text{length } a$   
**assumes**  $n < \text{length } b$   
**shows**  $(\text{map2 } f \ a \ b)!\ n = f \ (a!\ n) \ (b!\ n)$   
 ⟨*proof*⟩

Translating a polynomial by a vector.

**fun** *poly-translate* :: 'a list  $\Rightarrow$  'a poly  $\Rightarrow$  'a poly  
**where** *poly-translate*  $vs \ (poly.C \ c) = poly.C \ c$   
 | *poly-translate*  $vs \ (poly.Bound \ n) = poly.Add \ (poly.Bound \ n) \ (poly.C \ (vs \ ! \ n))$   
 | *poly-translate*  $vs \ (poly.Add \ l \ r) = poly.Add \ (poly-translate \ vs \ l) \ (poly-translate \ vs \ r)$   
 | *poly-translate*  $vs \ (poly.Sub \ l \ r) = poly.Sub \ (poly-translate \ vs \ l) \ (poly-translate \ vs \ r)$   
 | *poly-translate*  $vs \ (poly.Mul \ l \ r) = poly.Mul \ (poly-translate \ vs \ l) \ (poly-translate \ vs \ r)$   
 | *poly-translate*  $vs \ (poly.Neg \ p) = poly.Neg \ (poly-translate \ vs \ p)$   
 | *poly-translate*  $vs \ (poly.Pw \ p \ n) = poly.Pw \ (poly-translate \ vs \ p) \ n$   
 | *poly-translate*  $vs \ (poly.CN \ c \ n \ p) = poly.Add \ (poly-translate \ vs \ c) \ (poly.Mul \ (poly.Add \ (poly.Bound \ n) \ (poly.C \ (vs \ ! \ n))) \ (poly-translate \ vs \ p))$

Translating a polynomial is equivalent to translating its argument.

**lemma** *poly-translate-correct*:  
**assumes**  $\text{num-params } p \leq \text{length } x$   
**assumes**  $\text{length } x = \text{length } v$   
**shows**  $Ipoly \ x \ (poly-translate \ v \ p) = Ipoly \ (\text{map2 } (+) \ x \ v) \ p$   
 ⟨*proof*⟩

**lemma** *real-poly-translate*:  
**assumes**  $\text{num-params } p \leq \text{length } v$   
**shows**  $Ipoly \ x \ (\text{map-poly } \text{real-of-float} \ (poly-translate \ v \ p)) = Ipoly \ x \ (poly-translate \ v \ (\text{map-poly } \text{real-of-float} \ p))$   
 ⟨*proof*⟩

**lemma** *num-params-poly-translate[simp]*:  
**shows**  $\text{num-params } (poly-translate \ v \ p) = \text{num-params } p$   
 ⟨*proof*⟩

**end**

**theory** *Taylor-Models-Misc*

**imports**

*HOL-Library.Float*

*HOL-Library.Function-Algebras*

*HOL-Decision-Procs.Approximation*

*Affine-Arithmetic.Floatarith-Expression*

**begin**

This theory contains anything that doesn't belong anywhere else.

**lemma** *of-nat-real-float-equiv*:  $(\text{of-nat } n :: \text{real}) = (\text{of-nat } n :: \text{float})$   
 ⟨proof⟩

**lemma** *fact-real-float-equiv*:  $(\text{fact } n :: \text{float}) = (\text{fact } n :: \text{real})$   
 ⟨proof⟩

**lemma** *Some-those-length*:  
 $\text{those } ys = \text{Some } xs \implies \text{length } xs = \text{length } ys$   
 ⟨proof⟩

**lemma** *those-eq-None-iff*:  $\text{those } ys = \text{None} \iff \text{None} \in \text{set } ys$   
 ⟨proof⟩

**lemma** *those-eq-Some-iff*:  $\text{those } ys = (\text{Some } xs) \iff (ys = \text{map } \text{Some } xs)$   
 ⟨proof⟩

**lemma** *Some-those-nth*:  
**assumes**  $\text{those } ys = \text{Some } xs$   
**assumes**  $i < \text{length } xs$   
**shows**  $\text{Some } (xs!i) = ys!i$   
 ⟨proof⟩

**lemma** *fun-pow*:  $f^n = (\lambda x. (f x)^n)$   
 ⟨proof⟩

**context includes** *floatarith-notation* **begin**

Translate floatarith expressions by a vector of floats.

**fun** *fa-translate* ::  $\text{float list} \Rightarrow \text{floatarith} \Rightarrow \text{floatarith}$   
**where** *fa-translate*  $v$   $(\text{Add } a b) = \text{Add } (\text{fa-translate } v a) (\text{fa-translate } v b)$   
 | *fa-translate*  $v$   $(\text{Minus } a) = \text{Minus } (\text{fa-translate } v a)$   
 | *fa-translate*  $v$   $(\text{Mult } a b) = \text{Mult } (\text{fa-translate } v a) (\text{fa-translate } v b)$   
 | *fa-translate*  $v$   $(\text{Inverse } a) = \text{Inverse } (\text{fa-translate } v a)$   
 | *fa-translate*  $v$   $(\text{Cos } a) = \text{Cos } (\text{fa-translate } v a)$   
 | *fa-translate*  $v$   $(\text{Arctan } a) = \text{Arctan } (\text{fa-translate } v a)$   
 | *fa-translate*  $v$   $(\text{Min } a b) = \text{Min } (\text{fa-translate } v a) (\text{fa-translate } v b)$   
 | *fa-translate*  $v$   $(\text{Max } a b) = \text{Max } (\text{fa-translate } v a) (\text{fa-translate } v b)$   
 | *fa-translate*  $v$   $(\text{Abs } a) = \text{Abs } (\text{fa-translate } v a)$   
 | *fa-translate*  $v$   $(\text{Sqrt } a) = \text{Sqrt } (\text{fa-translate } v a)$   
 | *fa-translate*  $v$   $(\text{Exp } a) = \text{Exp } (\text{fa-translate } v a)$   
 | *fa-translate*  $v$   $(\text{Ln } a) = \text{Ln } (\text{fa-translate } v a)$   
 | *fa-translate*  $v$   $(\text{Var } n) = \text{Add } (\text{Var } n) (\text{Num } (v!n))$   
 | *fa-translate*  $v$   $(\text{Power } a n) = \text{Power } (\text{fa-translate } v a) n$   
 | *fa-translate*  $v$   $(\text{Powr } a b) = \text{Powr } (\text{fa-translate } v a) (\text{fa-translate } v b)$   
 | *fa-translate*  $v$   $(\text{Floor } x) = \text{Floor } (\text{fa-translate } v x)$   
 | *fa-translate*  $v$   $(\text{Num } c) = \text{Num } c$   
 | *fa-translate*  $v$   $\text{Pi} = \text{Pi}$

**lemma** *fa-translate-correct*:  
**assumes** *max-Var-floatarith f ≤ length I*  
**assumes** *length v = length I*  
**shows** *interpret-floatarith (fa-translate v f) I = interpret-floatarith f (map2 (+) I v)*  
 ⟨*proof*⟩

**primrec** *vars-floatarith* **where**

*vars-floatarith (Add a b) = (vars-floatarith a) ∪ (vars-floatarith b)*  
 | *vars-floatarith (Mult a b) = (vars-floatarith a) ∪ (vars-floatarith b)*  
 | *vars-floatarith (Inverse a) = vars-floatarith a*  
 | *vars-floatarith (Minus a) = vars-floatarith a*  
 | *vars-floatarith (Num a) = {}*  
 | *vars-floatarith (Var i) = {i}*  
 | *vars-floatarith (Cos a) = vars-floatarith a*  
 | *vars-floatarith (Arctan a) = vars-floatarith a*  
 | *vars-floatarith (Abs a) = vars-floatarith a*  
 | *vars-floatarith (Max a b) = (vars-floatarith a) ∪ (vars-floatarith b)*  
 | *vars-floatarith (Min a b) = (vars-floatarith a) ∪ (vars-floatarith b)*  
 | *vars-floatarith (Pi) = {}*  
 | *vars-floatarith (Sqrt a) = vars-floatarith a*  
 | *vars-floatarith (Exp a) = vars-floatarith a*  
 | *vars-floatarith (Powr a b) = (vars-floatarith a) ∪ (vars-floatarith b)*  
 | *vars-floatarith (Ln a) = vars-floatarith a*  
 | *vars-floatarith (Power a n) = vars-floatarith a*  
 | *vars-floatarith (Floor a) = vars-floatarith a*

**lemma** *finite-vars-floatarith[simp]*: *finite (vars-floatarith x)*  
 ⟨*proof*⟩

**end**

**lemma** *max-Var-floatarith-eq-Max-vars-floatarith*:

*max-Var-floatarith fa = (if vars-floatarith fa = {} then 0 else Suc (Max (vars-floatarith fa)))*  
 ⟨*proof*⟩

**end**

**theory** *Taylor-Models*

**imports** *Interval-Approximation*

*Horner-Eval*

*Polynomial-Expression-Additional*

*Taylor-Models-Misc*

*HOL-Decision-Proc.Approximation*

*HOL-Library.Function-Algebras*

*HOL-Library.Set-Algebras*

*Affine-Arithmetic.Straight-Line-Program*

*Affine-Arithmetic.Affine-Approximation*

**begin**

TODO: get rid of float poly/float interval and use real poly/real interval and data refinement?

## 7 Multivariate Taylor Models

### 7.1 Computing interval bounds on arithmetic expressions

This is a wrapper around the "approx" function. It computes range bounds on floatarith expressions.

```
fun compute-bound-fa :: nat ⇒ floatarith ⇒ float interval list ⇒ float interval
option
  where compute-bound-fa prec f I =
    (case approx prec f (map (Some o (λx. (lower x, upper x))) I) of
      Some (a, b) ⇒ (if a ≤ b then Some (Ivl a b) else None)
    | - ⇒ None)
```

```
lemma compute-bound-fa-correct:
  interpret floatarith f i ∈r ivl
  if compute-bound-fa prec f I = Some ivl
    i all-in I
  for i::real list
  ⟨proof⟩
```

### 7.2 Definition of Taylor models and notion of rangeity

Taylor models are a pair of a polynomial and an absolute error bound.

```
datatype taylor-model = TaylorModel (tm-poly: float poly) (tm-bound: float interval)
```

Taylor model for a real valuation of variables

```
primrec insertion :: (nat ⇒ 'a) ⇒ 'a poly ⇒ 'a::{plus,zero,minus,uminus,times,one,power}
where
  insertion bs (C c) = c
| insertion bs (poly.Bound n) = bs n
| insertion bs (Neg a) = - insertion bs a
| insertion bs (poly.Add a b) = insertion bs a + insertion bs b
| insertion bs (Sub a b) = insertion bs a - insertion bs b
| insertion bs (Mul a b) = insertion bs a * insertion bs b
| insertion bs (Pw t n) = insertion bs t ^ n
| insertion bs (CN c n p) = insertion bs c + (bs n) * insertion bs p
```

```
definition range-tm :: (nat ⇒ real) ⇒ taylor-model ⇒ real interval where
range-tm e tm = interval-of (insertion e (tm-poly tm)) + real-interval (tm-bound
tm)
```

```
lemma Ipoly-num-params-cong: Ipoly xs p = Ipoly ys p
  if ∧i. i < num-params p ⇒ xs ! i = ys ! i
```

*<proof>*

**lemma** *insertion-num-params-cong*: *insertion e p = insertion f p*  
**if**  $\bigwedge i. i < \text{num-params } p \implies e \ i = f \ i$   
*<proof>*

**lemma** *insertion-eq-IpolyI*: *insertion xs p = Ipoly ys p*  
**if**  $\bigwedge i. i < \text{num-params } p \implies xs \ i = ys \ ! \ i$   
*<proof>*

**lemma** *Ipoly-eq-insertionI*: *Ipoly ys p = insertion xs p*  
**if**  $\bigwedge i. i < \text{num-params } p \implies xs \ i = ys \ ! \ i$   
*<proof>*

**lemma** *range-tmI*:  
 $x \in_i \text{range-tm } e \ tm$   
**if**  $x: x \in_i \text{interval-of } (\text{insertion } e \ ((\text{tm-poly } tm))) + \text{real-interval } (\text{tm-bound } tm)$   
**for**  $e::\text{nat} \Rightarrow \text{real}$   
*<proof>*

**lemma** *range-tmD*:  
 $x \in_i \text{interval-of } (\text{insertion } e \ (\text{tm-poly } tm)) + \text{real-interval } (\text{tm-bound } tm)$   
**if**  $x \in_i \text{range-tm } e \ tm$   
**for**  $e::\text{nat} \Rightarrow \text{real}$   
*<proof>*

### 7.3 Interval bounds for Taylor models

Bound a polynomial by simply approximating it with interval arguments.

**fun** *compute-bound-poly* ::  $\text{nat} \Rightarrow \text{float interval poly} \Rightarrow (\text{float interval list}) \Rightarrow (\text{float interval list}) \Rightarrow \text{float interval}$  **where**  
  *compute-bound-poly prec (poly.C f) I a = f*  
  | *compute-bound-poly prec (poly.Bound n) I a = round-interval prec (I ! n - (a ! n))*  
  | *compute-bound-poly prec (poly.Add p q) I a =*  
    *round-interval prec (compute-bound-poly prec p I a + compute-bound-poly prec q I a)*  
  | *compute-bound-poly prec (poly.Sub p q) I a =*  
    *round-interval prec (compute-bound-poly prec p I a - compute-bound-poly prec q I a)*  
  | *compute-bound-poly prec (poly.Mul p q) I a =*  
    *mult-float-interval prec (compute-bound-poly prec p I a) (compute-bound-poly prec q I a)*  
  | *compute-bound-poly prec (poly.Neg p) I a = -compute-bound-poly prec p I a*  
  | *compute-bound-poly prec (poly.Pw p n) I a = power-float-interval prec n (compute-bound-poly prec p I a)*  
  | *compute-bound-poly prec (poly.CN p n q) I a =*  
    *round-interval prec (compute-bound-poly prec p I a +*  
      *mult-float-interval prec (round-interval prec (I ! n - (a ! n))) (compute-bound-poly*

$prec\ q\ I\ a))$

Bounds on Taylor models are simply a bound on its polynomial, widened by the approximation error.

**fun** *compute-bound-tm* ::  $nat \Rightarrow float\ interval\ list \Rightarrow float\ interval\ list \Rightarrow taylor\ model \Rightarrow float\ interval$   
**where** *compute-bound-tm*  $prec\ I\ a\ (TaylorModel\ p\ e) = compute-bound-poly\ prec\ p\ I\ a + e$

**lemma** *compute-bound-tm-def*:

$compute-bound-tm\ prec\ I\ a\ tm = compute-bound-poly\ prec\ (tm-poly\ tm)\ I\ a + (tm-bound\ tm)$   
(*proof*)

**lemma** *real-of-float-in-real-interval-of*[*intro, simp*]:  $real-of-float\ x \in_r X$  **if**  $x \in_i X$   
(*proof*)

**lemma** *in-set-of-round-interval*[*intro, simp*]:

$x \in_r round-interval\ prec\ X$  **if**  $x \in_r X$   
(*proof*)

**lemma** *in-set-real-minus-interval*[*intro, simp*]:

$x - y \in_r X - Y$  **if**  $x \in_r X\ y \in_r Y$   
(*proof*)

**lemma** *real-interval-plus*:  $real-interval\ (a + b) = real-interval\ a + real-interval\ b$   
(*proof*)

**lemma** *real-interval-uminus*:  $real-interval\ (- b) = - real-interval\ b$   
(*proof*)

**lemma** *real-interval-of*:  $real-interval\ (interval-of\ b) = interval-of\ b$   
(*proof*)

**lemma** *real-interval-minus*:  $real-interval\ (a - b) = real-interval\ a - real-interval\ b$   
(*proof*)

**lemma** *in-set-real-plus-interval*[*intro, simp*]:

$x + y \in_r X + Y$  **if**  $x \in_r X\ y \in_r Y$   
(*proof*)

**lemma** *in-set-neg-plus-interval*[*intro, simp*]:

$- y \in_r - Y$  **if**  $y \in_r Y$   
(*proof*)

**lemma** *real-interval-times*:  $real-interval\ (a * b) = real-interval\ a * real-interval\ b$   
(*proof*)

**lemma** *in-set-real-times-interval*[*intro, simp*]:

$x * y \in_r X * Y$  **if**  $x \in_r X$   $y \in_r Y$

*<proof>*

**lemma** *real-interval-one*: *real-interval* 1 = 1

*<proof>*

**lemma** *real-interval-zero*: *real-interval* 0 = 0

*<proof>*

**lemma** *real-interval-power*: *real-interval* (a ^ b) = *real-interval* a ^ b

*<proof>*

**lemma** *in-set-real-power-interval*[*intro, simp*]:

$x ^ n \in_r X ^ n$  **if**  $x \in_r X$

*<proof>*

**lemma** *power-float-interval-real-interval*[*intro, simp*]:

$x ^ n \in_r$  *power-float-interval* prec n X **if**  $x \in_r X$

*<proof>*

**lemma** *in-set-mult-float-interval*[*intro, simp*]:

$x * y \in_r$  *mult-float-interval* prec X Y **if**  $x \in_r X$   $y \in_r Y$

*<proof>*

**lemma** *in-set-real-minus-swap*I:  $e \in_r I \rightarrow i - a \in_r I$

**if**  $x - e \in_r a \in_r I \rightarrow x \in_r I$

*<proof>*

**definition** *develops-at-within*::(nat  $\Rightarrow$  real)  $\Rightarrow$  float interval list  $\Rightarrow$  float interval list  $\Rightarrow$  bool

**where** *develops-at-within* e a I  $\longleftrightarrow$  (a all-subset I)  $\wedge$  ( $\forall i < \text{length } I. e \in_r I \rightarrow i - a \in_r I$ )

**lemma** *develops-at-withinI*:

**assumes** all-in: a all-subset I

**assumes** e:  $\bigwedge i. i < \text{length } I \Rightarrow e \in_r I \rightarrow i - a \in_r I$

**shows** *develops-at-within* e a I

*<proof>*

**lemma** *develops-at-withinD*:

**assumes** *develops-at-within* e a I

**shows** a all-subset I

$\bigwedge i. i < \text{length } I \Rightarrow e \in_r I \rightarrow i - a \in_r I$

*<proof>*

**lemma** *compute-bound-poly-correct*:

**fixes** p::float poly

**assumes** num-params p  $\leq$  length I

**assumes** *dev*: *develops-at-within e a I*  
**shows** *insertion e (p::real poly) ∈<sub>r</sub> compute-bound-poly prec (map-poly interval-of p) I a*  
 ⟨*proof*⟩

**lemma** *compute-bound-tm-correct*:  
**fixes** *I :: float interval list and f :: real list ⇒ real*  
**assumes** *n: num-params (tm-poly t) ≤ length I*  
**assumes** *dev: develops-at-within e a I*  
**assumes** *x0: x0 ∈<sub>i</sub> range-tm e t*  
**shows** *x0 ∈<sub>r</sub> compute-bound-tm prec I a t*  
 ⟨*proof*⟩

**lemma** *compute-bound-tm-correct-subset*:  
**fixes** *I :: float interval list and f :: real list ⇒ real*  
**assumes** *n: num-params (tm-poly t) ≤ length I*  
**assumes** *dev: develops-at-within e a I*  
**shows** *set-of (range-tm e t) ⊆ set-of (real-interval (compute-bound-tm prec I a t))*  
 ⟨*proof*⟩

**lemma** *compute-bound-poly-mono*:  
**assumes** *num-params p ≤ length I*  
**assumes** *mem: I all-subset J a all-subset I*  
**shows** *set-of (compute-bound-poly prec p I a) ⊆ set-of (compute-bound-poly prec p J a)*  
 ⟨*proof*⟩

**lemma** *compute-bound-tm-mono*:  
**fixes** *I :: float interval list and f :: real list ⇒ real*  
**assumes** *num-params (tm-poly t) ≤ length I*  
**assumes** *I all-subset J*  
**assumes** *a all-subset I*  
**shows** *set-of (compute-bound-tm prec I a t) ⊆ set-of (compute-bound-tm prec J a t)*  
 ⟨*proof*⟩

## 7.4 Computing taylor models for basic, univariate functions

**definition** *tm-const :: float ⇒ taylor-model*  
**where** *tm-const c = TaylorModel (poly.C c) 0*

**context includes** *floatarith-notation* **begin**

**definition** *tm-pi :: nat ⇒ taylor-model*  
**where** *tm-pi prec = (*  
*let pi-ivl = the (compute-bound-fa prec Pi [])*  
*in TaylorModel (poly.C (mid pi-ivl)) (centered pi-ivl)*  
*)*

**lemma** *zero-real-interval*[*intro,simp*]:  $0 \in_r 0$   
 ⟨*proof*⟩

**lemma** *range-TM-tm-const*[*simp*]:  $\text{range-tm } e \text{ (tm-const } c) = \text{interval-of } c$   
 ⟨*proof*⟩

**lemma** *num-params-tm-const*[*simp*]:  $\text{num-params (tm-poly (tm-const } c)) = 0$   
 ⟨*proof*⟩

**lemma** *num-params-tm-pi*[*simp*]:  $\text{num-params (tm-poly (tm-pi } prec)) = 0$   
 ⟨*proof*⟩

**lemma** *range-tm-tm-pi*:  $pi \in_i \text{range-tm } e \text{ (tm-pi } prec)$   
 ⟨*proof*⟩

### 7.4.1 Derivations of floatarith expressions

Compute the  $n$ th derivative of a floatarith expression

**fun** *deriv* ::  $\text{nat} \Rightarrow \text{floatarith} \Rightarrow \text{nat} \Rightarrow \text{floatarith}$   
**where**  $\text{deriv } v \ f \ 0 = f$   
 |  $\text{deriv } v \ f \ (\text{Suc } n) = \text{DERIV-floatarith } v \ (\text{deriv } v \ f \ n)$

**lemma** *isDERIV-DERIV-floatarith*:  
**assumes**  $\text{isDERIV } v \ f \ vs$   
**shows**  $\text{isDERIV } v \ (\text{DERIV-floatarith } v \ f) \ vs$   
 ⟨*proof*⟩

**lemma** *isDERIV-is-analytic*:  
 $\text{isDERIV } i \ (\text{Taylor-Models.deriv } i \ f \ n) \ xs$   
**if**  $\text{isDERIV } i \ f \ xs$   
 ⟨*proof*⟩

**lemma** *deriv-correct*:  
**assumes**  $\text{isDERIV } i \ f \ (xs[i:=t]) \ i < \text{length } xs$   
**shows**  $((\lambda x. \text{interpret-floatarith } (\text{deriv } i \ f \ n) \ (xs[i:=x])) \text{ has-real-derivative } \text{interpret-floatarith } (\text{deriv } i \ f \ (\text{Suc } n)) \ (xs[i:=t]))$   
 (at  $t$  within  $S$ )  
 ⟨*proof*⟩

Faster derivation for univariate functions, producing smaller terms and thus less over-approximation.

TODO: Extend to Arctan, Log!

**fun** *deriv-rec* ::  $\text{floatarith} \Rightarrow \text{nat} \Rightarrow \text{floatarith}$   
**where**  $\text{deriv-rec } (\text{Exp } (\text{Var } 0)) \ n = \text{Exp } (\text{Var } 0)$   
 |  $\text{deriv-rec } (\text{Cos } (\text{Var } 0)) \ n = (\text{case } n \ \text{mod } 4$   
   of  $0 \Rightarrow \text{Cos } (\text{Var } 0)$   
   |  $\text{Suc } 0 \Rightarrow \text{Minus } (\text{Sin } (\text{Var } 0))$

```

| Suc (Suc 0) ⇒ Minus (Cos (Var 0))
| Suc (Suc (Suc 0)) ⇒ Sin (Var 0))
| deriv-rec (Inverse (Var 0)) n = (if n = 0 then Inverse (Var 0) else Mult (Num
(fact n * (if n mod 2 = 0 then 1 else -1))) (Inverse (Power (Var 0) (Suc n))))
| deriv-rec f n = deriv 0 f n

```

**lemma** *deriv-rec-correct*:

```

assumes isDERIV 0 f (xs[0:=t]) 0 < length xs
shows ((λx. interpret-floatarith (deriv-rec f n) (xs[0:=x])) has-real-derivative
interpret-floatarith (deriv-rec f (Suc n)) (xs[0:=t])) (at t within S)
⟨proof⟩

```

**lemma** *deriv-rec-0-idem[simp]*:

```

shows deriv-rec f 0 = f
⟨proof⟩

```

## 7.4.2 Computing Taylor models for arbitrary univariate expressions

```

fun tmf-c :: nat ⇒ float interval list ⇒ floatarith ⇒ nat ⇒ float interval option
where tmf-c prec I f i = compute-bound-fa prec (Mult (deriv-rec f i) (Inverse
(Num (fact i)))) I

```

— The interval coefficients of the Taylor polynomial, i.e. the real coefficients approximated by a float interval.

```

fun tmf-ivl-cs :: nat ⇒ nat ⇒ float interval list ⇒ float list ⇒ floatarith ⇒ float
interval list option

```

```

where tmf-ivl-cs prec ord I a f = those (map (tmf-c prec a f) [0..<ord] @ [tmf-c
prec I f ord])

```

— Make a list of bounds on the n+1 coefficients, with the n+1-th coefficient bounding the remainder term of the Taylor-Lagrange formula.

```

fun tmf-polys :: float interval list ⇒ float poly × float interval poly

```

```

where tmf-polys [] = (poly.C 0, poly.C 0)
| tmf-polys (c # cs) = (
  let (pf, pi) = tmf-polys cs
  in (poly.CN (poly.C (mid c)) 0 pf, poly.CN (poly.C (centered c)) 0 pi)
)

```

```

fun tm-floatarith :: nat ⇒ nat ⇒ float interval list ⇒ float list ⇒ floatarith ⇒
taylor-model option

```

```

where tm-floatarith prec ord I a f = (
  map-option (λcs.
    let (pf, pi) = tmf-polys cs;
    - = compute-bound-tm prec (List.map2 (−) I a);
    e = round-interval prec (Ipoly (List.map2 (−) I a) pi) — TODO: use
compute-bound-tm here?!
    in TaylorModel pf e
  ) (tmf-ivl-cs prec ord I a f)

```

) — Compute a Taylor model from an arbitrary, univariate floatarith expression, if possible. This is used to compute Taylor models for elemental functions like sin, cos, exp, etc.

**term** *compute-bound-poly*

**lemma** *tmf-c-correct*:

**fixes** *A::float interval list and I::float interval and f::floatarith and a::real list*

**assumes** *a all-in A*

**assumes** *tmf-c prec A f i = Some I*

**shows** *interpret-floatarith (deriv-rec f i) a / fact i ∈<sub>r</sub> I*

*<proof>*

**lemma** *tmf-ivl-cs-length*:

**assumes** *tmf-ivl-cs prec n A a f = Some cs*

**shows** *length cs = n + 1*

*<proof>*

**lemma** *tmf-ivl-cs-correct*:

**fixes** *A::float interval list and f::floatarith*

**assumes** *a all-in I*

**assumes** *tmf-ivl-cs prec ord I a f = Some cs*

**shows**  $\bigwedge i. i < \text{ord} \implies \text{tmf-c prec (map interval-of a) f i} = \text{Some (cs!i)}$

**and** *tmf-c prec I f ord = Some (cs!ord)*

**and** *length cs = Suc ord*

*<proof>*

**lemma** *Ipoly-fst-tmf-polys*:

*Ipoly xs (fst (tmf-polys z)) = ( $\sum i < \text{length } z. xs ! 0 \wedge i * (\text{mid } (z ! i))$ )*

**for** *xs::real list*

*<proof>*

**lemma** *insertion-fst-tmf-polys*:

*insertion e (fst (tmf-polys z)) = ( $\sum i < \text{length } z. e 0 \wedge i * (\text{mid } (z ! i))$ )*

**for** *e::nat  $\Rightarrow$  real*

*<proof>*

**lemma** *Ipoly-snd-tmf-polys*:

*set-of (horner-eval (real-interval o centered o nth z) x (length z))  $\subseteq$  set-of (Ipoly [x] (map-poly real-interval (snd (tmf-polys z))))*

*<proof>*

**lemma** *zero-interval[*intro,simp*]: 0 ∈<sub>i</sub> 0*

*<proof>*

**lemma** *sum-in-intervalI: sum f X ∈<sub>i</sub> sum g X if  $\bigwedge x. x \in X \implies f x \in_i g x$*

**for** *f :: -  $\Rightarrow$  'a :: ordered-comm-monoid-add*

*<proof>*

**lemma** *set-of-sum-subset: set-of (sum f X)  $\subseteq$  set-of (sum g X)*

**if**  $\bigwedge x. x \in X \implies \text{set-of } (f x) \subseteq \text{set-of } (g x)$   
**for**  $f :: \text{--}'a::\text{linordered-ab-group-add interval}$   
 <proof>

**lemma interval-of-plus:**  $\text{interval-of } (a + b) = \text{interval-of } a + \text{interval-of } b$   
 <proof>

**lemma interval-of-uminus:**  $\text{interval-of } (- a) = - \text{interval-of } a$   
 <proof>

**lemma interval-of-zero:**  $\text{interval-of } 0 = 0$   
 <proof>

**lemma interval-of-sum:**  $\text{interval-of } (\text{sum } f X) = \text{sum } (\lambda x. \text{interval-of } (f x)) X$   
 <proof>

**lemma interval-of-prod:**  $\text{interval-of } (a * b) = \text{interval-of } a * \text{interval-of } b$   
 <proof>

**lemma in-set-of-interval-of[simp]:**  $x \in_i (\text{interval-of } y) \longleftrightarrow x = y$  **for**  $x y::'a::\text{order}$   
 <proof>

**lemma real-interval-Ipoly:**  $\text{real-interval } (Ipoly xs p) = Ipoly (\text{map real-interval } xs)$   
 ( $\text{map-poly real-interval } p$ )  
**if**  $\text{num-params } p \leq \text{length } xs$   
 <proof>

**lemma num-params-tmf-polys1:**  $\text{num-params } (\text{fst } (\text{tmf-polys } z)) \leq \text{Suc } 0$   
 <proof>

**lemma num-params-tmf-polys2:**  $\text{num-params } (\text{snd } (\text{tmf-polys } z)) \leq \text{Suc } 0$   
 <proof>

**lemma set-of-real-interval-subset:**  $\text{set-of } (\text{real-interval } x) \subseteq \text{set-of } (\text{real-interval } y)$   
**if**  $\text{set-of } x \subseteq \text{set-of } y$   
 <proof>

**theorem tm-floatarith:**  
**assumes**  $t: \text{tm-floatarith prec ord } I xs f = \text{Some } t$   
**assumes**  $a: xs \text{ all-in } I$  **and**  $x: x \in_r I ! 0$   
**assumes**  $xs\text{-ne}: xs \neq []$   
**assumes**  $\text{deriv}: \bigwedge x. x \in_r I ! 0 \implies \text{isDERIV } 0 f (xs[0 := x])$   
**assumes**  $\bigwedge i. 0 < i \implies i < \text{length } xs \implies e i = \text{real-of-float } (xs ! i)$   
**assumes**  $\text{diff-e}: (x - \text{real-of-float } (xs ! 0)) = e 0$   
**shows**  $\text{interpret-floatarith } f (xs[0:=x]) \in_i \text{range-tm } e t$   
 <proof>

## 7.5 Operations on Taylor models

**fun** *tm-norm-poly* :: *taylor-model*  $\Rightarrow$  *taylor-model*

**where** *tm-norm-poly* (*TaylorModel* *p e*) = *TaylorModel* (*polynat* *p*) *e*

— Normalizes the Taylor model by transforming its polynomial into horner form.

**fun** *tm-lower-order* *tm-lower-order-of-normed* :: *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *float interval list*  $\Rightarrow$  *float interval list*  $\Rightarrow$  *taylor-model*  $\Rightarrow$  *taylor-model*

**where** *tm-lower-order* *prec ord I a t* = *tm-lower-order-of-normed* *prec ord I a* (*tm-norm-poly* *t*)

| *tm-lower-order-of-normed* *prec ord I a* (*TaylorModel* *p e*) = (  
     *let* (*l, r*) = *split-by-degree* *ord p*  
     *in TaylorModel* *l* (*round-interval* *prec* (*e* + *compute-bound-poly* *prec r I a*))  
 )

— Reduces the degree of a Taylor model's polynomial to *n* and keeps it range by increasing the error bound.

**fun** *tm-round-floats* *tm-round-floats-of-normed* :: *nat*  $\Rightarrow$  *float interval list*  $\Rightarrow$  *float interval list*  $\Rightarrow$  *taylor-model*  $\Rightarrow$  *taylor-model*

**where** *tm-round-floats* *prec I a t* = *tm-round-floats-of-normed* *prec I a* (*tm-norm-poly* *t*)

| *tm-round-floats-of-normed* *prec I a* (*TaylorModel* *p e*) = (  
     *let* (*l, r*) = *split-by-prec* *prec p*  
     *in TaylorModel* *l* (*round-interval* *prec* (*e* + *compute-bound-poly* *prec r I a*))  
 )

— Rounding of Taylor models. Rounds both the coefficients of the polynomial and the floats in the error bound.

**fun** *tm-norm* *tm-norm'* :: *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *float interval list*  $\Rightarrow$  *float interval list*  $\Rightarrow$  *taylor-model*  $\Rightarrow$  *taylor-model*

**where** *tm-norm* *prec ord I a t* = *tm-norm'* *prec ord I a* (*tm-norm-poly* *t*)

| *tm-norm'* *prec ord I a t* = *tm-round-floats-of-normed* *prec I a* (*tm-lower-order-of-normed* *prec ord I a t*)

— Normalization of taylor models. Performs order lowering and rounding on taylor models, also converts the polynomial into horner form.

**fun** *tm-neg* :: *taylor-model*  $\Rightarrow$  *taylor-model*

**where** *tm-neg* (*TaylorModel* *p e*) = *TaylorModel* ( $\sim_p$  *p*) ( $-e$ )

**fun** *tm-add* :: *taylor-model*  $\Rightarrow$  *taylor-model*  $\Rightarrow$  *taylor-model*

**where** *tm-add* (*TaylorModel* *p1 e1*) (*TaylorModel* *p2 e2*) = *TaylorModel* (*p1* +<sub>*p*</sub> *p2*) (*e1* + *e2*)

**fun** *tm-sub* :: *taylor-model*  $\Rightarrow$  *taylor-model*  $\Rightarrow$  *taylor-model*

**where** *tm-sub* *t1 t2* = *tm-add* *t1* (*tm-neg* *t2*)

**fun** *tm-mul* :: *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *float interval list*  $\Rightarrow$  *float interval list*  $\Rightarrow$  *taylor-model*  $\Rightarrow$  *taylor-model*  $\Rightarrow$  *taylor-model*

**where** *tm-mul* *prec ord I a* (*TaylorModel* *p1 e1*) (*TaylorModel* *p2 e2*) = (  
     *let* *d1* = *compute-bound-poly* *prec p1 I a*;

```

      d2 = compute-bound-poly prec p2 I a;
      p = p1 *_p p2;
      e = e1*d2 + d1*e2 + e1*e2
    in tm-norm' prec ord I a (TaylorModel p e)
  )
lemmas [simp del] = tm-norm'.simps

fun tm-pow :: nat ⇒ nat ⇒ float interval list ⇒ float interval list ⇒ taylor-model
⇒ nat ⇒ taylor-model
  where tm-pow prec ord I a t 0 = tm-const 1
  | tm-pow prec ord I a t (Suc n) = (
    if odd (Suc n)
    then tm-mul prec ord I a t (tm-pow prec ord I a t n)
    else let t' = tm-pow prec ord I a t ((Suc n) div 2)
         in tm-mul prec ord I a t' t'
  )

```

Evaluates a float polynomial, using a Taylor model as the parameter. This is used to compose Taylor models.

```

fun eval-poly-at-tm :: nat ⇒ nat ⇒ float interval list ⇒ float interval list ⇒ float
poly ⇒ taylor-model ⇒ taylor-model
  where eval-poly-at-tm prec ord I a (poly.C c) t = tm-const c
  | eval-poly-at-tm prec ord I a (poly.Bound n) t = t
  | eval-poly-at-tm prec ord I a (poly.Add p1 p2) t
    = tm-add (eval-poly-at-tm prec ord I a p1 t)
              (eval-poly-at-tm prec ord I a p2 t)
  | eval-poly-at-tm prec ord I a (poly.Sub p1 p2) t
    = tm-sub (eval-poly-at-tm prec ord I a p1 t)
              (eval-poly-at-tm prec ord I a p2 t)
  | eval-poly-at-tm prec ord I a (poly.Mul p1 p2) t
    = tm-mul prec ord I a (eval-poly-at-tm prec ord I a p1 t)
                          (eval-poly-at-tm prec ord I a p2 t)
  | eval-poly-at-tm prec ord I a (poly.Neg p) t
    = tm-neg (eval-poly-at-tm prec ord I a p t)
  | eval-poly-at-tm prec ord I a (poly.Pw p n) t
    = tm-pow prec ord I a (eval-poly-at-tm prec ord I a p t) n
  | eval-poly-at-tm prec ord I a (poly.CN c n p) t = (
    let pt = eval-poly-at-tm prec ord I a p t;
        t-mul-pt = tm-mul prec ord I a t pt
    in tm-add (eval-poly-at-tm prec ord I a c t) t-mul-pt
  )

```

```

fun tm-inc-err :: float interval ⇒ taylor-model ⇒ taylor-model
  where tm-inc-err i (TaylorModel p e) = TaylorModel p (e + i)

```

```

fun tm-comp :: nat ⇒ nat ⇒ float interval list ⇒ float interval list ⇒ float ⇒
taylor-model ⇒ taylor-model ⇒ taylor-model
  where tm-comp prec ord I a ta (TaylorModel p e) t = (
    let t-sub-ta = tm-sub t (tm-const ta);

```

```

    pt = eval-poly-at-tm prec ord I a p t-sub-ta
  in tm-inc-err e pt
)

```

*tm-max*, *tm-min* and *tm-abs* are implemented extremely naively, because I don't expect them to be very useful. But the implementation is fairly modular, i.e. *tm*-{*abs,min,max*} all can easily be swapped out, as long as the corresponding correctness lemmas *tm*-{*abs,min,max*}-*range* are updated as well.

```

fun tm-abs :: nat => float interval list => float interval list => taylor-model =>
taylor-model

```

```

  where tm-abs prec I a t = (
    let bound = compute-bound-tm prec I a t; abs-bound=Ivl (0::float) (max (abs
(lower bound)) (abs (upper bound)))
    in TaylorModel (poly.C (mid abs-bound)) (centered abs-bound))

```

```

fun tm-union :: nat => float interval list => float interval list => taylor-model =>
taylor-model => taylor-model

```

```

  where tm-union prec I a t1 t2 = (
    let b1 = compute-bound-tm prec I a t1; b2 = compute-bound-tm prec I a t2;
    b-combined = sup b1 b2
    in TaylorModel (poly.C (mid b-combined)) (centered b-combined))

```

```

fun tm-min :: nat => float interval list => float interval list => taylor-model =>
taylor-model => taylor-model

```

```

  where tm-min prec I a t1 t2 = tm-union prec I a t1 t2

```

```

fun tm-max :: nat => float interval list => float interval list => taylor-model =>
taylor-model => taylor-model

```

```

  where tm-max prec I a t1 t2 = tm-union prec I a t1 t2

```

Rangeity of is preserved by our operations on Taylor models.

```

lemma insertion-polyadd[simp]: insertion e (a +p b) = insertion e a + insertion
e b

```

```

  for a b::'a::ring-1 poly
  <proof>

```

```

lemma insertion-polyneq[simp]: insertion e (~p b) = - insertion e b

```

```

  for b::'a::ring-1 poly
  <proof>

```

```

lemma insertion-polysub[simp]: insertion e (a -p b) = insertion e a - insertion
e b

```

```

  for a b::'a::ring-1 poly
  <proof>

```

```

lemma insertion-polymul[simp]: insertion e (a *p b) = insertion e a * insertion e
b

```

**for**  $a\ b :: 'a :: \text{comm-ring-1 poly}$   
 <proof>

**lemma** *insertion-polypow[simp]*:  $\text{insertion } e (a \hat{^}_p b) = \text{insertion } e a \hat{^} b$   
**for**  $a :: 'a :: \text{comm-ring-1 poly}$   
 <proof>

**lemma** *insertion-polynate [simp]*:  
 $\text{insertion } bs (\text{polynate } p) = (\text{insertion } bs p :: 'a :: \text{comm-ring-1})$   
 <proof>

**lemma** *tm-norm-poly-range*:  
**assumes**  $x \in_i \text{range-tm } e\ t$   
**shows**  $x \in_i \text{range-tm } e (tm\text{-norm-poly } t)$   
 <proof>

**lemma** *split-by-degree-correct-insertion*:  
**fixes**  $x :: \text{nat} \Rightarrow \text{real}$  **and**  $p :: \text{float poly}$   
**assumes** *split-by-degree*  $\text{ord } p = (l, r)$   
**shows**  $\text{maxdegree } l \leq \text{ord}$  (**is** ?P1)  
**and**  $\text{insertion } x\ p = \text{insertion } x\ l + \text{insertion } x\ r$  (**is** ?P2)  
**and**  $\text{num-params } l \leq \text{num-params } p$  (**is** ?P3)  
**and**  $\text{num-params } r \leq \text{num-params } p$  (**is** ?P4)  
 <proof>

**lemma** *split-by-prec-correct-insertion*:  
**fixes**  $x :: \text{nat} \Rightarrow \text{real}$  **and**  $p :: \text{float poly}$   
**assumes** *split-by-prec*  $\text{ord } p = (l, r)$   
**shows**  $\text{insertion } x\ p = \text{insertion } x\ l + \text{insertion } x\ r$  (**is** ?P1)  
**and**  $\text{num-params } l \leq \text{num-params } p$  (**is** ?P2)  
**and**  $\text{num-params } r \leq \text{num-params } p$  (**is** ?P3)  
 <proof>

**lemma** *tm-lower-order-of-normed-range*:  
**assumes**  $x \in_i \text{range-tm } e\ t$   
**assumes** *dev: develops-at-within*  $e\ a\ I$   
**assumes**  $\text{num-params } (tm\text{-poly } t) \leq \text{length } I$   
**shows**  $x \in_i \text{range-tm } e (tm\text{-lower-order-of-normed } prec\ \text{ord } I\ a\ t)$   
 <proof>

**lemma** *num-params-tm-norm-poly-le*:  $\text{num-params } (tm\text{-poly } (tm\text{-norm-poly } t)) \leq X$   
**if**  $\text{num-params } (tm\text{-poly } t) \leq X$   
 <proof>

**lemma** *tm-lower-order-range*:  
**assumes**  $x \in_i \text{range-tm } e\ t$   
**assumes** *dev: develops-at-within*  $e\ a\ I$   
**assumes**  $\text{num-params } (tm\text{-poly } t) \leq \text{length } I$

**shows**  $x \in_i \text{range-tm } e \text{ (tm-lower-order prec ord I a t)}$   
<proof>

**lemma** *tm-round-floats-of-normed-range:*

**assumes**  $x \in_i \text{range-tm } e \text{ t}$

**assumes** *dev: develops-at-within e a I*

**assumes**  $\text{num-params (tm-poly t)} \leq \text{length I}$

**shows**  $x \in_i \text{range-tm } e \text{ (tm-round-floats-of-normed prec I a t)}$

— TODO: this is a clone of  $[[?x \in_i \text{range-tm } ?e \text{ ?t}; \text{develops-at-within } ?e \text{ ?a } ?I;$   
 $\text{num-params (tm-poly ?t)} \leq \text{length ?I}] \implies ?x \in_i \text{range-tm } ?e \text{ (tm-lower-order-of-normed}$   
 $?prec \text{ ?ord } ?I \text{ ?a } ?t) \text{ -i general sweeping method!}$   
<proof>

**lemma** *num-params-split-by-degree-le:*  $\text{num-params (fst (split-by-degree ord x))} \leq K$

$\text{num-params (snd (split-by-degree ord x))} \leq K$

**if**  $\text{num-params } x \leq K$  **for**  $x::\text{float poly}$

<proof>

**lemma** *num-params-split-by-prec-le:*  $\text{num-params (fst (split-by-prec ord x))} \leq K$

$\text{num-params (snd (split-by-prec ord x))} \leq K$

**if**  $\text{num-params } x \leq K$  **for**  $x::\text{float poly}$

<proof>

**lemma** *num-params-tm-norm'-le:*

$\text{num-params (tm-poly (tm-round-floats-of-normed prec I a t))} \leq X$

**if**  $\text{num-params (tm-poly t)} \leq X$

<proof>

**lemma** *tm-round-floats-range:*

**assumes**  $x \in_i \text{range-tm } e \text{ t develops-at-within e a I num-params (tm-poly t)} \leq \text{length I}$

**shows**  $x \in_i \text{range-tm } e \text{ (tm-round-floats prec I a t)}$

<proof>

**lemma** *num-params-tm-lower-order-of-normed-le:*  $\text{num-params (tm-poly (tm-lower-order-of-normed prec ord I a t))} \leq X$

**if**  $\text{num-params (tm-poly t)} \leq X$

<proof>

**lemma** *tm-norm'-range:*

**assumes**  $x \in_i \text{range-tm } e \text{ t develops-at-within e a I num-params (tm-poly t)} \leq \text{length I}$

**shows**  $x \in_i \text{range-tm } e \text{ (tm-norm' prec ord I a t)}$

<proof>

**lemma** *num-params-tm-norm':*

$\text{num-params (tm-poly (tm-norm' prec ord I a t))} \leq X$

**if**  $\text{num-params } (tm\text{-poly } t) \leq X$   
 $\langle \text{proof} \rangle$

**lemma** *tm-norm-range*:

**assumes**  $x \in_i \text{range-tm } e \ t \ \text{develops-at-within } e \ a \ I \ \text{num-params } (tm\text{-poly } t) \leq$   
 $\text{length } I$

**shows**  $x \in_i \text{range-tm } e \ (tm\text{-norm } \text{prec } \text{ord } I \ a \ t)$   
 $\langle \text{proof} \rangle$

**lemmas**  $[\text{simp } \text{del}] = tm\text{-norm.simps}$

**lemma** *tm-neg-range*:

**assumes**  $x \in_i \text{range-tm } e \ t$

**shows**  $-x \in_i \text{range-tm } e \ (tm\text{-neg } t)$   
 $\langle \text{proof} \rangle$

**lemmas**  $[\text{simp } \text{del}] = tm\text{-neg.simps}$

**lemma** *tm-bound-tm-add[simp]*:  $tm\text{-bound } (tm\text{-add } t1 \ t2) = tm\text{-bound } t1 + tm\text{-bound } t2$

$\langle \text{proof} \rangle$

**lemma** *interval-of-add*:  $\text{interval-of } (a + b) = \text{interval-of } a + \text{interval-of } b$

$\langle \text{proof} \rangle$

**lemma** *tm-add-range*:

$x + y \in_i \text{range-tm } e \ (tm\text{-add } t1 \ t2)$

**if**  $x \in_i \text{range-tm } e \ t1$

$y \in_i \text{range-tm } e \ t2$

$\langle \text{proof} \rangle$

**lemmas**  $[\text{simp } \text{del}] = tm\text{-add.simps}$

**lemma** *tm-sub-range*:

**assumes**  $x \in_i \text{range-tm } e \ t1$

**assumes**  $y \in_i \text{range-tm } e \ t2$

**shows**  $x - y \in_i \text{range-tm } e \ (tm\text{-sub } t1 \ t2)$

$\langle \text{proof} \rangle$

**lemmas**  $[\text{simp } \text{del}] = tm\text{-sub.simps}$

**lemma** *set-of-intervalI*:  $\text{set-of } (\text{interval-of } y) \subseteq \text{set-of } Y$  **if**  $y \in_i Y$  **for**  $y::'a::\text{order}$

$\langle \text{proof} \rangle$

**lemma** *set-of-real-intervalI*:  $\text{set-of } (\text{interval-of } y) \subseteq \text{set-of } (\text{real-interval } Y)$  **if**  $y$

$\in_r Y$

$\langle \text{proof} \rangle$

**lemma** *tm-mul-range*:

**assumes**  $x \in_i \text{range-tm } e \ t1$

**assumes**  $y \in_i \text{range-tm } e \ t2$

**assumes**  $\text{dev}: \text{develops-at-within } e \ a \ I$

**assumes** *params*:  $\text{num-params } (tm\text{-poly } t1) \leq \text{length } I \text{ num-params } (tm\text{-poly } t2)$   
 $\leq \text{length } I$   
**shows**  $x * y \in_i \text{range-tm } e (tm\text{-mul } \text{prec } \text{ord } I \text{ a } t1 \ t2)$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-tm-mul-le*:  
 $\text{num-params } (tm\text{-poly } (tm\text{-mul } \text{prec } \text{ord } I \text{ a } t1 \ t2)) \leq X$   
**if**  $\text{num-params } (tm\text{-poly } t1) \leq X$   
 $\text{num-params } (tm\text{-poly } t2) \leq X$   
 $\langle \text{proof} \rangle$

**lemmas** [*simp del*] = *tm-pow.simps*— TODO: make a systematic decision

**lemma**  
**shows** *tm-pow-range*:  $\text{num-params } (tm\text{-poly } t) \leq \text{length } I \implies$   
 $\text{develops-at-within } e \text{ a } I \implies$   
 $x \in_i \text{range-tm } e \ t \implies$   
 $x \wedge n \in_i \text{range-tm } e (tm\text{-pow } \text{prec } \text{ord } I \text{ a } t \ n)$   
**and** *num-params-tm-pow-le*[*THEN order-trans*]:  
 $\text{num-params } (tm\text{-poly } (tm\text{-pow } \text{prec } \text{ord } I \text{ a } t \ n)) \leq \text{num-params } (tm\text{-poly } t)$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-tm-add-le*:  
 $\text{num-params } (tm\text{-poly } (tm\text{-add } t1 \ t2)) \leq X$   
**if**  $\text{num-params } (tm\text{-poly } t1) \leq X$   
 $\text{num-params } (tm\text{-poly } t2) \leq X$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-tm-neg-eq*[*simp*]:  
 $\text{num-params } (tm\text{-poly } (tm\text{-neg } t1)) = \text{num-params } (tm\text{-poly } t1)$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-tm-sub-le*:  
 $\text{num-params } (tm\text{-poly } (tm\text{-sub } t1 \ t2)) \leq X$   
**if**  $\text{num-params } (tm\text{-poly } t1) \leq X$   
 $\text{num-params } (tm\text{-poly } t2) \leq X$   
 $\langle \text{proof} \rangle$

**lemma** *num-params-eval-poly-le*:  $\text{num-params } (tm\text{-poly } (eval\text{-poly-at-tm } \text{prec } \text{ord } I \text{ a } p \ t)) \leq x$   
**if**  $\text{num-params } (tm\text{-poly } t) \leq x \text{ num-params } p \leq \max 1 \ x$   
 $\langle \text{proof} \rangle$

**lemma** *eval-poly-at-tm-range*:  
**assumes**  $\text{num-params } p \leq 1$   
**assumes** *tg-def*:  $e' 0 \in_i \text{range-tm } e \ tg$   
**assumes** *dev*:  $\text{develops-at-within } e \text{ a } I$  **and** *params*:  $\text{num-params } (tm\text{-poly } tg) \leq$   
 $\text{length } I$   
**shows** *insertion*  $e' p \in_i \text{range-tm } e (eval\text{-poly-at-tm } \text{prec } \text{ord } I \text{ a } p \ tg)$

*<proof>*

**lemma** *tm-inc-err-range*:  $x \in_i \text{range-tm } e \text{ (tm-inc-err } i \text{ } t)$   
**if**  $x \in_i \text{range-tm } e \text{ } t + \text{real-interval } i$   
*<proof>*

**lemma** *num-params-tm-inc-err*:  $\text{num-params (tm-poly (tm-inc-err } i \text{ } t))} \leq X$   
**if**  $\text{num-params (tm-poly } t) \leq X$   
*<proof>*

**lemma** *num-params-tm-comp-le*:  $\text{num-params (tm-poly (tm-comp prec ord } I \text{ } a \text{ } ga \text{ } tf \text{ } tg))} \leq X$   
**if**  $\text{num-params (tm-poly } tf) \leq \max 1 X \text{ num-params (tm-poly } tg) \leq X$   
*<proof>*

**lemma** *tm-comp-range*:  
**assumes** *tf-def*:  $x \in_i \text{range-tm } e' \text{ } tf$   
**assumes** *tg-def*:  $e' 0 \in_i \text{range-tm } e \text{ (tm-sub } tg \text{ (tm-const } ga))$   
**assumes** *params*:  $\text{num-params (tm-poly } tf) \leq 1 \text{ num-params (tm-poly } tg) \leq$   
*length } I*  
**assumes** *dev*: *develops-at-within } e a I*  
**shows**  $x \in_i \text{range-tm } e \text{ (tm-comp prec ord } I \text{ } a \text{ } ga \text{ } tf \text{ } tg)$   
*<proof>*

**lemma** *mid-centered-collapse*:  
 $\text{interval-of (real-of-float (mid abs-bound))} + \text{real-interval (centered abs-bound)} =$   
 $\text{real-interval abs-bound}$   
*<proof>*

**lemmas** [*simp del*] = *tm-abs.simps*

**lemma** *tm-abs-range*:  
**assumes** *x*:  $x \in_i \text{range-tm } e \text{ } t$   
**assumes** *n*:  $\text{num-params (tm-poly } t) \leq \text{length } I$  **and** *d*: *develops-at-within } e a I*  
**shows**  $\text{abs } x \in_i \text{range-tm } e \text{ (tm-abs prec } I \text{ } a \text{ } t)$   
*<proof>*

**lemma** *num-params-tm-abs-le*:  $\text{num-params (tm-poly (tm-abs prec } I \text{ } a \text{ } t))} \leq X$  **if**  
 $\text{num-params (tm-poly } t) \leq X$   
*<proof>*

**lemma** *real-interval-sup*:  $\text{real-interval (sup } a \text{ } b) = \text{sup (real-interval } a) \text{ (real-interval } b)$   
*<proof>*

**lemma** *in-interval-supI1*:  $x \in_i a \implies x \in_i \text{sup } a \text{ } b$   
**and** *in-interval-supI2*:  $x \in_i b \implies x \in_i \text{sup } a \text{ } b$   
**for**  $x :: 'a :: \text{lattice}$   
*<proof>*

**lemma** *tm-union-range-left*:  
**assumes**  $x \in_i \text{range-tm } e \ t1$   
 $\text{num-params } (tm\text{-poly } t1) \leq \text{length } I \text{ develops-at-within } e \ a \ I$   
**shows**  $x \in_i \text{range-tm } e \ (tm\text{-union } prec \ I \ a \ t1 \ t2)$   
 $\langle proof \rangle$

**lemma** *tm-union-range-right*:  
**assumes**  $x \in_i \text{range-tm } e \ t2$   
 $\text{num-params } (tm\text{-poly } t2) \leq \text{length } I \text{ develops-at-within } e \ a \ I$   
**shows**  $x \in_i \text{range-tm } e \ (tm\text{-union } prec \ I \ a \ t1 \ t2)$   
 $\langle proof \rangle$

**lemma** *num-params-tm-union-le*:  
 $\text{num-params } (tm\text{-poly } (tm\text{-union } prec \ I \ a \ t1 \ t2)) \leq X$   
**if**  $\text{num-params } (tm\text{-poly } t1) \leq X \ \text{num-params } (tm\text{-poly } t2) \leq X$   
 $\langle proof \rangle$

**lemmas** [*simp del*] = *tm-union.simps tm-min.simps tm-max.simps*

**lemma** *tm-min-range*:  
**assumes**  $x \in_i \text{range-tm } e \ t1$   
**assumes**  $y \in_i \text{range-tm } e \ t2$   
 $\text{num-params } (tm\text{-poly } t1) \leq \text{length } I$   
 $\text{num-params } (tm\text{-poly } t2) \leq \text{length } I$   
 $\text{develops-at-within } e \ a \ I$   
**shows**  $\min \ x \ y \in_i \text{range-tm } e \ (tm\text{-min } prec \ I \ a \ t1 \ t2)$   
 $\langle proof \rangle$

**lemma** *tm-max-range*:  
**assumes**  $x \in_i \text{range-tm } e \ t1$   
**assumes**  $y \in_i \text{range-tm } e \ t2$   
 $\text{num-params } (tm\text{-poly } t1) \leq \text{length } I$   
 $\text{num-params } (tm\text{-poly } t2) \leq \text{length } I$   
 $\text{develops-at-within } e \ a \ I$   
**shows**  $\max \ x \ y \in_i \text{range-tm } e \ (tm\text{-max } prec \ I \ a \ t1 \ t2)$   
 $\langle proof \rangle$

## 7.6 Computing Taylor models for multivariate expressions

Compute Taylor models for expressions of the form "f (g x)", where f is an elementary function like exp or cos, by composing Taylor models for f and g. For our correctness proof, we need to make it explicit that the range of g on I is inside the domain of f, by introducing the *f-exists-on* predicate.

**fun** *compute-tm-by-comp* ::  $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{float interval list} \Rightarrow \text{float interval list} \Rightarrow$   
 $\text{floatarith} \Rightarrow \text{taylor-model option} \Rightarrow (\text{float interval} \Rightarrow \text{bool}) \Rightarrow \text{taylor-model option}$   
**where** *compute-tm-by-comp* *prec ord I a f g f-exists-on* = (  
 $\text{case } g$   
 $\text{of } \text{Some } tg \Rightarrow ($

```

    let gI = compute-bound-tm prec I a tg;
        ga = mid (compute-bound-tm prec a a tg)
    in if f-exists-on gI
        then map-option ( $\lambda$ tf. tm-comp prec ord I a ga tf tg ) (tm-floatarith
prec ord [gI] [ga] f)
        else None)
    | -  $\Rightarrow$  None
)

```

Compute Taylor models with numerical precision  $prec$  of degree  $ord$ , with Taylor models in the environment  $env$  whose variables are jointly interpreted with domain  $I$  and expanded around point  $a$ . from floatarith expressions on a rectangular domain.

```

fun approx-tm :: nat  $\Rightarrow$  nat  $\Rightarrow$  float interval list  $\Rightarrow$  float interval list  $\Rightarrow$  floatarith
 $\Rightarrow$  taylor-model list  $\Rightarrow$ 
    taylor-model option
where approx-tm - - I - (Num c) env = Some (tm-const c)
| approx-tm - - I a (Var n) env = (if n < length env then Some (env ! n) else
None)
| approx-tm prec ord I a (Add l r) env = (
    case (approx-tm prec ord I a l env, approx-tm prec ord I a r env)
    of (Some t1, Some t2)  $\Rightarrow$  Some (tm-add t1 t2)
    | -  $\Rightarrow$  None)
| approx-tm prec ord I a (Minus f) env
    = map-option tm-neg (approx-tm prec ord I a f env)
| approx-tm prec ord I a (Mult l r) env = (
    case (approx-tm prec ord I a l env, approx-tm prec ord I a r env)
    of (Some t1, Some t2)  $\Rightarrow$  Some (tm-mul prec ord I a t1 t2)
    | -  $\Rightarrow$  None)
| approx-tm prec ord I a (Power f k) env
    = map-option ( $\lambda$ t. tm-pow prec ord I a t k)
        (approx-tm prec ord I a f env)
| approx-tm prec ord I a (Inverse f) env
    = compute-tm-by-comp prec ord I a (Inverse (Var 0)) (approx-tm prec ord
I a f env) ( $\lambda$ x. 0 < lower x  $\vee$  upper x < 0)
| approx-tm prec ord I a (Cos f) env
    = compute-tm-by-comp prec ord I a (Cos (Var 0)) (approx-tm prec ord I a
f env) ( $\lambda$ x. True)
| approx-tm prec ord I a (Arctan f) env
    = compute-tm-by-comp prec ord I a (Arctan (Var 0)) (approx-tm prec ord
I a f env) ( $\lambda$ x. True)
| approx-tm prec ord I a (Exp f) env
    = compute-tm-by-comp prec ord I a (Exp (Var 0)) (approx-tm prec ord I a
f env) ( $\lambda$ x. True)
| approx-tm prec ord I a (Ln f) env
    = compute-tm-by-comp prec ord I a (Ln (Var 0)) (approx-tm prec ord I a f
env) ( $\lambda$ x. 0 < lower x)
| approx-tm prec ord I a (Sqrt f) env
    = compute-tm-by-comp prec ord I a (Sqrt (Var 0)) (approx-tm prec ord I a

```

$f \text{ env}$  ( $\lambda x. 0 < \text{lower } x$ )  
 |  $\text{approx-tm prec ord } I \text{ a } Pi \text{ env} = \text{Some } (tm\text{-pi prec})$   
 |  $\text{approx-tm prec ord } I \text{ a } (Abs f) \text{ env}$   
    $= \text{map-option } (tm\text{-abs prec } I \text{ a}) (\text{approx-tm prec ord } I \text{ a } f \text{ env})$   
 |  $\text{approx-tm prec ord } I \text{ a } (Min l r) \text{ env} =$   
    $\text{case } (\text{approx-tm prec ord } I \text{ a } l \text{ env}, \text{approx-tm prec ord } I \text{ a } r \text{ env})$   
    $\text{of } (Some t1, Some t2) \Rightarrow \text{Some } (tm\text{-min prec } I \text{ a } t1 t2)$   
   |  $- \Rightarrow \text{None}$   
 |  $\text{approx-tm prec ord } I \text{ a } (Max l r) \text{ env} =$   
    $\text{case } (\text{approx-tm prec ord } I \text{ a } l \text{ env}, \text{approx-tm prec ord } I \text{ a } r \text{ env})$   
    $\text{of } (Some t1, Some t2) \Rightarrow \text{Some } (tm\text{-max prec } I \text{ a } t1 t2)$   
   |  $- \Rightarrow \text{None}$   
 |  $\text{approx-tm prec ord } I \text{ a } (Powr l r) \text{ env} = \text{None} \text{ — TODO}$   
 |  $\text{approx-tm prec ord } I \text{ a } (Floor l) \text{ env} = \text{None} \text{ — TODO}$

**lemma** *mid-in-real-interval*:  $\text{mid } i \in_r i$   
 $\langle \text{proof} \rangle$

**lemma** *set-of-real-interval-mono*:  $\text{set-of } (\text{real-interval } x) \subseteq \text{set-of } (\text{real-interval } y)$   
**if**  $\text{set-of } x \subseteq \text{set-of } y$   
 $\langle \text{proof} \rangle$

**lemmas** [*simp del*] = *compute-bound-poly.simps tm-floatarith.simps*

**lemmas** [*simp del*] = *tmf-ivl-cs.simps compute-bound-tm.simps tmf-polys.simps*

**lemma** *tm-floatarith-eq-Some-num-params*:  
 $\text{tm-floatarith prec ord } a \text{ b } f = \text{Some } tf \implies \text{num-params } (tm\text{-poly } tf) \leq 1$   
 $\langle \text{proof} \rangle$

**lemma** *compute-tm-by-comp-range*:  
**assumes**  $\text{max-Var-floatarith } f \leq 1$   
**assumes**  $a: a \text{ all-subset } I$   
**assumes**  $tx\text{-range}: x \in_i \text{range-tm } e \text{ tg}$   
**assumes**  $t\text{-def}: \text{compute-tm-by-comp prec ord } I \text{ a } f (\text{Some } tg) \text{ c} = \text{Some } t$   
**assumes**  $f\text{-deriv}$ :  
 $\bigwedge x. x \in_r \text{compute-bound-tm prec } I \text{ a } tg \implies c (\text{compute-bound-tm prec } I \text{ a } tg)$   
 $\implies \text{isDERIV } 0 \text{ f } [x]$   
**assumes**  $\text{params}: \text{num-params } (tm\text{-poly } tg) \leq \text{length } I$   
**and**  $\text{dev}: \text{develops-at-within } e \text{ a } I$   
**shows**  $\text{interpret-floatarith } f [x] \in_i \text{range-tm } e \text{ t}$   
 $\langle \text{proof} \rangle$

**lemmas** [*simp del*] = *compute-tm-by-comp.simps*

**lemma** *compute-tm-by-comp-num-params-le*:  
**assumes**  $\text{compute-tm-by-comp prec ord } I \text{ a } f (\text{Some } t0) \text{ i} = \text{Some } t$

**assumes**  $1 \leq X$  *num-params* (tm-poly t0)  $\leq X$   
**shows** *num-params* (tm-poly t)  $\leq X$   
 ⟨proof⟩

**lemma** *compute-tm-by-comp-eq-Some-iff*: *compute-tm-by-comp prec ord I a f t0 i*  
 = *Some t*  $\longleftrightarrow$   
 ( $\exists z x2. t0 = \text{Some } x2 \wedge$   
   *tm-floatarith prec ord [compute-bound-tm prec I a x2]*  
   [*mid (compute-bound-tm prec a a x2)*] *f* =  
   *Some z*  
 $\wedge$  *tm-comp prec ord I a*  
   (*mid (compute-bound-tm prec a a x2)*)  $z x2 = t$   
 $\wedge i$  (*compute-bound-tm prec I a x2*))  
 ⟨proof⟩

**lemma** *num-params-approx-tm*:  
**assumes** *approx-tm prec ord I a f env = Some t*  
**assumes**  $\bigwedge tm. tm \in \text{set } env \implies \text{num-params } (tm\text{-poly } tm) \leq \text{length } I$   
**shows** *num-params (tm-poly t)  $\leq$  length I*  
 ⟨proof⟩

**lemma** *in-interval-reall*:  $a \in_i I$  **if**  $a \in_r I$  ⟨proof⟩

**lemma** *all-subset-all-inI*: *map interval-of a all-subset I* **if** *a all-in I*  
 ⟨proof⟩

**lemma** *compute-tm-by-comp-None*: *compute-tm-by-comp p ord I a x None k =*  
*None*  
 ⟨proof⟩

**lemma** *approx-tm-num-Vars-None*:  
**assumes** *max-Var-floatarith f > length env*  
**shows** *approx-tm p ord I a f env = None*  
 ⟨proof⟩

**lemma** *approx-tm-num-Vars*:  
**assumes** *approx-tm prec ord I a f env = Some t*  
**shows** *max-Var-floatarith f  $\leq$  length env*  
 ⟨proof⟩

**definition** *range-tms e xs = map (range-tm e) xs*

**lemma** *approx-tm-range*:  
**assumes** *a: a all-subset I*  
**assumes** *t-def: approx-tm prec ord I a f env = Some t*  
**assumes** *allin: xs all-in<sub>i</sub> range-tms e env*  
**assumes** *devs: develops-at-within e a I*  
**assumes** *env:  $\bigwedge tm. tm \in \text{set } env \implies \text{num-params } (tm\text{-poly } tm) \leq \text{length } I$*   
**shows** *interpret-floatarith f xs  $\in_i$  range-tm e t*

*<proof>*

Evaluate expression with Taylor models in environment.

## 7.7 Computing bounds for floatarith expressions

TODO: compare parametrization of input vs. uncertainty for input...

**definition**  $tm\text{-of-ivl-par } n \text{ ivl} = TaylorModel (CN (C ((upper \text{ ivl} + lower \text{ ivl}) * Float 1 (-1))) n$   
 $(C ((upper \text{ ivl} - lower \text{ ivl}) * Float 1 (-1)))) 0$   
— track uncertainty in parameter  $n$ , which is to be interpreted over standardized domain  $[-1, 1]$ .

**value**  $tm\text{-of-ivl-par } 3 \text{ (Ivl } (-1) 1)$

**definition**  $tms\text{-of-ivls } ivls = map (\lambda(i, ivl). tm\text{-of-ivl-par } i \text{ ivl}) (zip [0..<length ivls] ivls)$

**value**  $tms\text{-of-ivls } [Ivl 1 2, Ivl 4 5]$

**primrec**  $approx\text{-slp}'::nat \Rightarrow nat \Rightarrow float \text{ interval list} \Rightarrow float \text{ interval list} \Rightarrow slp$   
 $\Rightarrow$

$taylor\text{-model list} \Rightarrow taylor\text{-model list option}$

**where**

$approx\text{-slp}' p \text{ ord } I a [] \text{ xs} = Some \text{ xs}$   
 $| approx\text{-slp}' p \text{ ord } I a (ea \# eas) \text{ xs} =$   
   $do \{$   
     $r \leftarrow approx\text{-tm } p \text{ ord } I a ea \text{ xs};$   
     $approx\text{-slp}' p \text{ ord } I a eas (r \# xs)$   
   $\}$

**lemma**  $mem\text{-range-tms-Cons-iff}[simp]: x \# xs \text{ all-in}_i \text{ range-tms } e (X \# XS) \longleftrightarrow x$   
 $\in_i \text{ range-tm } e X \wedge xs \text{ all-in}_i \text{ range-tms } e XS$   
*<proof>*

**lemma**  $approx\text{-slp}'\text{-range}:$

**assumes**  $i: i \text{ all-subset } I$

**assumes**  $dev: \text{ develops-at-within } e i I$

**assumes**  $vs: vs \text{ all-in}_i \text{ range-tms } e VS (\wedge tm. tm \in \text{ set } VS \implies \text{ num-params}$   
 $(tm\text{-poly } tm) \leq \text{ length } I)$

**assumes**  $appr: approx\text{-slp}' p \text{ ord } I i \text{ ra } VS = Some X$

**shows**  $interpret\text{-slp } ra \text{ vs all-in}_i \text{ range-tms } e X$

*<proof>*

**definition**  $approx\text{-slp}::nat \Rightarrow nat \Rightarrow nat \Rightarrow slp \Rightarrow taylor\text{-model list} \Rightarrow taylor\text{-model}$   
 $list \text{ option}$

**where**

$approx\text{-slp } p \text{ ord } d \text{ slp } tms =$   
 $map\text{-option } (take d)$

(*approx-slp' p ord (replicate (length tms) (Ivl (-1) 1)) (replicate (length tms) 0) slp tms*)

**lemma** *length-range-tms[simp]*: *length (range-tms e VS) = length VS*  
 ⟨*proof*⟩

**lemma** *set-of-Ivl*: *set-of (Ivl a b) = {a .. b} if a ≤ b*  
 ⟨*proof*⟩

**lemma** *set-of-zero[simp]*: *set-of 0 = {0::'a::ordered-comm-monoid-add}*  
 ⟨*proof*⟩

**theorem** *approx-slp-range-tms*:

**assumes** *approx-slp p ord d slp VS = Some X*

**assumes** *slp-def: slp = slp-of-fas fas*

**assumes** *d-def: d = length fas*

**assumes** *e: e ∈ UNIV → {-1 .. 1}*

**assumes** *vs: vs all-in<sub>i</sub> range-tms e VS*

**assumes** *lens: ∧tm. tm ∈ set VS ⇒ num-params (tm-poly tm) ≤ length vs*

**shows** *interpret-floatariths fas vs all-in<sub>i</sub> range-tms e X*

⟨*proof*⟩

**end**

**end**

**theory** *Experiments*

**imports** *Taylor-Models*

*Affine-Arithmetic.Affine-Arithmetic*

**begin**

**instantiation** *interval::({show, preorder}) show begin*

**lift-definition** *shows-prec-interval::*

*nat ⇒ 'a interval ⇒ char list ⇒ char list*

**is** *λp ivl s. (shows-string "Interval" o shows ivl) s* ⟨*proof*⟩

**lift-definition** *shows-list-interval::*

*'a interval list ⇒ char list ⇒ char list*

**is** *λivls s. shows-list ivls s* ⟨*proof*⟩

**instance**

⟨*proof*⟩

**end**

**definition** *split-largest-interval :: float interval list ⇒ float interval list × float interval list* **where**

*split-largest-interval xs = (case sort-key (uminus o snd) (zip [0..<length xs] (map*

( $\lambda x. \text{upper } x - \text{lower } x$ )  $xs$ ) of Nil  $\Rightarrow$  ([], [])  
 | ( $i, -$ )#-  $\Rightarrow$  let  $x = xs! i$  in ( $xs[i:=Ivl (\text{lower } x) ((\text{upper } x + \text{lower } x)*\text{Float } 1 (-1))]$ ),  
 $xs[i:=Ivl ((\text{upper } x + \text{lower } x)*\text{Float } 1 (-1)) (\text{upper } x)]$ )

**definition** *Inf-tm p params tm =*

*lower (compute-bound-tm p (replicate params (Ivl (-1) (1))) (replicate params (Ivl 0 0)) tm)*

**primrec** *prove-pos::bool  $\Rightarrow$  nat  $\Rightarrow$  nat  $\Rightarrow$  nat  $\Rightarrow$*

*(nat  $\Rightarrow$  nat  $\Rightarrow$  taylor-model list  $\Rightarrow$  taylor-model option)  $\Rightarrow$  float interval list list  $\Rightarrow$  bool* **where**

*prove-pos prnt 0 p ord F X = (let - = if prnt then print (STR "# depth limit exceeded"  $\leftarrow$ )"') else () in False)*

| *prove-pos prnt (Suc i) p ord F XXS =*  
*(case XXS of []  $\Rightarrow$  True | (X#XS)  $\Rightarrow$  let*

*params = length X;*

*R = F p ord (tms-of-ivls X);*

*- = if prnt then print (String.implode ((shows "# " o shows (map ( $\lambda ivl. (\text{lower } ivl, \text{upper } ivl)) X$ )) " $\leftarrow$ "))'') else ()*

*in*

*if  $R \neq \text{None} \wedge 0 < \text{Inf-tm } p \text{ params (the } R)$*

*then let - = if prnt then print (STR "# Success"  $\leftarrow$ )"') else () in prove-pos prnt i p ord F XS*

*else let - = if prnt then print (String.implode ((shows "# Split (" o shows ((map-option (Inf-tm p params)) R) o shows "'")' " $\leftarrow$ "))'') else () in case split-largest-interval X of (a, b)  $\Rightarrow$*

*prove-pos prnt i p ord F (a#b#XS))*

**hide-const** (open) *prove-pos-slp*

**definition** *prove-pos-slp prnt prec ord fa i xs = (let slp = slp-of-fas [fa] in prove-pos prnt i prec ord ( $\lambda p$  ord xs.*

*case approx-slp prec ord 1 slp xs of None  $\Rightarrow$  None | Some [x]  $\Rightarrow$  Some x | Some -  $\Rightarrow$  None) xs)*

**experiment begin**

**unbundle** *floatarith-notation*

**abbreviation** *schwefel  $\equiv$*

*(5.8806 / 10 ^ 10) + (Var 0 - (Var 1) ^ e2) ^ e2 + (Var 1 - 1) ^ e2 + (Var 0 - (Var 2) ^ e2) ^ e2 + (Var 2 - 1) ^ e2*

**lemma** *prove-pos-slp True 30 0 schwefel 100000 [replicate 3 (Ivl (-10) 10)]*

*(proof)*

**abbreviation** *delta6  $\equiv$  (Var 0 \* Var 3 \* (-Var 0 + Var 1 + Var 2 - Var 3 +*

$$\begin{aligned}
& \text{Var } 4 + \text{Var } 5) + \\
& \quad \text{Var } 1 * \text{Var } 4 * (\text{Var } 0 - \text{Var } 1 + \text{Var } 2 + \text{Var } 3 - \text{Var } 4 + \text{Var } 5) + \\
& \quad \text{Var } 2 * \text{Var } 5 * (\text{Var } 0 + \text{Var } 1 - \text{Var } 2 + \text{Var } 3 + \text{Var } 4 - \text{Var } 5) \\
& - \text{Var } 1 * \text{Var } 2 * \text{Var } 3 \\
& - \text{Var } 0 * \text{Var } 2 * \text{Var } 4 \\
& - \text{Var } 0 * \text{Var } 1 * \text{Var } 5 \\
& - \text{Var } 3 * \text{Var } 4 * \text{Var } 5)
\end{aligned}$$

**lemma** *prove-pos-slp True 30 3 delta6 10000 [replicate 6 (Ivl 4 (Float 104045 (-14)))]*  
*<proof>*

**abbreviation** *caprasse*  $\equiv (3.1801 + - \text{Var } 0 * (\text{Var } 2) \hat{e} 3 + 4 * \text{Var } 1 * (\text{Var } 2) \hat{e} 2 * \text{Var } 3 +$   
 $4 * \text{Var } 0 * \text{Var } 2 * (\text{Var } 3) \hat{e} 2 + 2 * \text{Var } 1 * (\text{Var } 3) \hat{e} 3 + 4 * \text{Var } 0 * \text{Var } 2 + 4 * (\text{Var } 2) \hat{e} 2 - 10 * \text{Var } 1 * \text{Var } 3 +$   
 $-10 * (\text{Var } 3) \hat{e} 2 + 2)$

**lemma** *prove-pos-slp True 30 2 caprasse 10000 [replicate 4 (Ivl (-Float 1 (-1)) (Float 1 (-1)))]*  
*<proof>*

**abbreviation** *magnetism*  $\equiv$   
 $0.25001 + (\text{Var } 0) \hat{e} 2 + 2 * (\text{Var } 1) \hat{e} 2 + 2 * (\text{Var } 2) \hat{e} 2 + 2 * (\text{Var } 3) \hat{e} 2$   
 $+ 2 * (\text{Var } 4) \hat{e} 2 + 2 * (\text{Var } 5) \hat{e} 2 +$   
 $2 * (\text{Var } 6) \hat{e} 2 - \text{Var } 0$

**end**

**end**