

Standard Borel Spaces

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Abstract

This entry includes a formalization of standard Borel spaces and (a variant of) the Borel isomorphism theorem. A separable complete metrizable topological space is called a polish space and a measurable space generated from a polish space is called a standard Borel space. We formalize the notion of standard Borel spaces by establishing set-based metric spaces, and then prove (a variant of) the Borel isomorphism theorem. The theorem states that a standard Borel spaces is either a countable discrete space or isomorphic to \mathbb{R} .

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We refer to the HOL-Analysis library, the textbooks by Matsuzaka [2] and Srivastava [3], and the lecture note by Biskup [1].

1 Lemmas

```
theory Lemmas-StandardBorel
  imports HOL-Probability.Probability
begin
```

1.1 Lemmas for Abstract Topology

1.1.1 Generated By

```
lemma topology-generated-by-sub:
  assumes  $\bigwedge U. U \in \mathcal{U} \implies (\text{openin } X \ U)$ 
  and  $\text{openin } (\text{topology-generated-by } \mathcal{U}) \ U$ 
  shows  $\text{openin } X \ U$ 
proof -
  have generate-topology-on  $\mathcal{U} \ U$ 
  by (simp add: assms(2) openin-topology-generated-by)
  then show ?thesis
  by induction (use assms(1) in auto)
qed
```

```
lemma topology-generated-by-open:
   $S = \text{topology-generated-by } \{U \mid U . \text{openin } S \ U\}$ 
  unfolding topology-eq
proof standard+
  fix  $U$ 
  assume  $\text{openin } (\text{topology-generated-by } \{U \mid U . \text{openin } S \ U\}) \ U$ 
  note this[simplified openin-topology-generated-by-iff]
  then show  $\text{openin } S \ U$ 
```

by induction auto
qed(simp add: openin-topology-generated-by-iff generate-topology-on.Basis)

lemma topology-generated-by-eq:

assumes $\bigwedge U. U \in \mathcal{U} \implies (\text{openin } (\text{topology-generated-by } \mathcal{O}) U)$
and $\bigwedge U. U \in \mathcal{O} \implies (\text{openin } (\text{topology-generated-by } \mathcal{U}) U)$
shows $\text{topology-generated-by } \mathcal{O} = \text{topology-generated-by } \mathcal{U}$
using topology-generated-by-sub[of \mathcal{U} , OF assms(1)] topology-generated-by-sub[of \mathcal{O} , OF assms(2)]
by(auto simp: topology-eq)

lemma topology-generated-by-homeomorphic-spaces:

assumes homeomorphic-map $X Y f X = \text{topology-generated-by } \mathcal{O}$
shows $Y = \text{topology-generated-by } ((\cdot) f \cdot \mathcal{O})$
unfolding topology-eq

proof

have $f: \text{open-map } X Y f \text{ inj-on } f (\text{topspace } X)$
using assms(1) **by** (simp-all add: homeomorphic-imp-open-map perfect-injective-eq-homeomorphic-map[sym])
obtain g **where** $g: \bigwedge x. x \in \text{topspace } X \implies g(f x) = x \wedge y. y \in \text{topspace } Y \implies f(g y) = y$ **open-map** $Y X g \text{ inj-on } g (\text{topspace } Y)$
using homeomorphic-map-maps[of $X Y f$,simplified assms(1)] homeomorphic-imp-open-map
homeomorphic-maps-map[of $X Y f$] homeomorphic-imp-injective-map[of $Y X$] **by**
blast

show $\bigwedge S. \text{openin } Y S = \text{openin } (\text{topology-generated-by } ((\cdot) f \cdot \mathcal{O})) S$

proof safe

fix S

assume $\text{openin } Y S$

then have $\text{openin } X (g \cdot S)$

using $g(3)$ **by** (simp add: open-map-def)

hence $h: \text{generate-topology-on } \mathcal{O} (g \cdot S)$

by(simp add: assms(2) openin-topology-generated-by-iff)

have $S = f \cdot (g \cdot S)$

using openin-subset[OF $\langle \text{openin } Y S \rangle$] $g(2)$ **by**(fastforce simp: image-def)

also have $\text{openin } (\text{topology-generated-by } ((\cdot) f \cdot \mathcal{O})) \dots$

using h

proof induction

case Empty

then show ?case **by** simp

next

case (Int $a b$)

with inj-on-image-Int[OF $f(2)$,of $a b$] **show** ?case

by (metis assms(2) openin-Int openin-subset openin-topology-generated-by-iff)

next

case (UN K)

then show ?case

by(auto simp: image-Union)

next

case (Basis s)

then show ?case

```

    by(auto intro!: generate-topology-on.Basis simp: openin-topology-generated-by-iff)
  qed
  finally show openin (topology-generated-by (( $\cdot$ )  $f^{-1} \mathcal{O}$ ))  $S$  .
next
fix  $S$ 
assume openin (topology-generated-by (( $\cdot$ )  $f^{-1} \mathcal{O}$ ))  $S$ 
then have generate-topology-on (( $\cdot$ )  $f^{-1} \mathcal{O}$ )  $S$ 
  by(simp add: openin-topology-generated-by-iff)
thus openin  $Y S$ 
proof induction
  case (Basis  $s$ )
  then obtain  $U$  where  $u: U \in \mathcal{O}$   $s = f^{-1} U$  by auto
  then show ?case
    using assms(1) assms(2) homeomorphic-map-openness-eq topology-generated-by-Basis
by blast
  qed auto
  qed
qed

```

```

lemma open-map-generated-topo:
  assumes  $\bigwedge u. u \in U \implies \text{openin } S (f^{-1} u) \text{ inj-on } f (\text{topspace } (\text{topology-generated-by } U))$ 
  shows open-map (topology-generated-by  $U$ )  $S$   $f$ 
  unfolding open-map-def
proof safe
  fix  $u$ 
  assume openin (topology-generated-by  $U$ )  $u$ 
  then have generate-topology-on  $U$   $u$ 
    by(simp add: openin-topology-generated-by-iff)
  thus openin  $S (f^{-1} u)$ 
proof induction
  case (Int  $a$   $b$ )
  then have [simp]:  $f^{-1} (a \cap b) = f^{-1} a \cap f^{-1} b$ 
  by (meson assms(2) inj-on-image-Int openin-subset openin-topology-generated-by-iff)
  from Int show ?case by auto
qed (simp-all add: image-Union openin-clauses(3) assms)
qed

```

```

lemma subtopology-generated-by:
  subtopology (topology-generated-by  $\mathcal{O}$ )  $T = \text{topology-generated-by } \{T \cap U \mid U. U \in \mathcal{O}\}$ 
  unfolding topology-eq openin-subtopology openin-topology-generated-by-iff
proof safe
  fix  $A$ 
  assume generate-topology-on  $\mathcal{O}$   $A$ 
  then show generate-topology-on  $\{T \cap U \mid U. U \in \mathcal{O}\}$   $(A \cap T)$ 
proof induction
  case Empty
  then show ?case

```

```

    by (simp add: generate-topology-on.Empty)
next
  case (Int a b)
  moreover have  $a \cap b \cap T = (a \cap T) \cap (b \cap T)$  by auto
  ultimately show ?case
    by(auto intro!: generate-topology-on.Int)
next
  case (UN K)
  moreover have  $(\bigcup K \cap T) = (\bigcup \{k \cap T \mid k. k \in K\})$  by auto
  ultimately show ?case
    by(auto intro!: generate-topology-on.UN)
next
  case (Basis s)
  then show ?case
    by(auto intro!: generate-topology-on.Basis)
qed
next
fix A
assume generate-topology-on  $\{T \cap U \mid U. U \in \mathcal{O}\} A$ 
then show  $\exists L. \text{generate-topology-on } \mathcal{O} L \wedge A = L \cap T$ 
proof induction
  case Empty
  show ?case
    by(auto intro!: exI[where x={}] generate-topology-on.Empty)
next
  case ih:(Int a b)
  then obtain La Lb where
    generate-topology-on  $\mathcal{O} La a = La \cap T$  generate-topology-on  $\mathcal{O} Lb b = Lb \cap T$ 
    by auto
  thus ?case
    using ih by(auto intro!: exI[where x= $La \cap Lb$ ] generate-topology-on.Int)
next
  case ih:(UN K)
  then obtain L where
 $\bigwedge k. k \in K \implies \text{generate-topology-on } \mathcal{O} (L k) \bigwedge k. k \in K \implies k = (L k) \cap T$ 
    by metis
  thus ?case
    using ih by(auto intro!: exI[where x= $\bigcup k \in K. L k$ ] generate-topology-on.UN)
next
  case (Basis s)
  then show ?case
    using generate-topology-on.Basis by fastforce
qed
qed

lemma prod-topology-generated-by:
  topology-generated-by  $\{U \times V \mid U V. U \in \mathcal{O} \wedge V \in \mathcal{U}\} = \text{prod-topology}$ 
  (topology-generated-by  $\mathcal{O}$ ) (topology-generated-by  $\mathcal{U}$ )
  unfolding topology-eq

```

```

proof safe
  fix U
  assume h:openin (topology-generated-by {U × V | U V. U ∈ O ∧ V ∈ U}) U
  show openin (prod-topology (topology-generated-by O) (topology-generated-by U))
  U
  by(auto simp: openin-prod-Times-iff[of topology-generated-by O topology-generated-by
  U]
    intro!: topology-generated-by-Basis topology-generated-by-sub[OF - h])
next
  fix U
  assume openin (prod-topology (topology-generated-by O) (topology-generated-by
  U)) U
  then have ∀z∈U. ∃ V1 V2. openin (topology-generated-by O) V1 ∧ openin
  (topology-generated-by U) V2 ∧ fst z ∈ V1 ∧ snd z ∈ V2 ∧ V1 × V2 ⊆ U
  by(auto simp: openin-prod-topology-alt)
  hence ∃ V1. ∀z∈U. ∃ V2. openin (topology-generated-by O) (V1 z) ∧ openin
  (topology-generated-by U) V2 ∧ fst z ∈ (V1 z) ∧ snd z ∈ V2 ∧ (V1 z) × V2 ⊆ U
  by(rule bchoice)
  then obtain V1 where ∀z∈U. ∃ V2. openin (topology-generated-by O) (V1 z)
  ∧ openin (topology-generated-by U) V2 ∧ fst z ∈ (V1 z) ∧ snd z ∈ V2 ∧ (V1 z)
  × V2 ⊆ U
  by auto
  hence ∃ V2. ∀z∈U. openin (topology-generated-by O) (V1 z) ∧ openin (topology-generated-by
  U) (V2 z) ∧ fst z ∈ (V1 z) ∧ snd z ∈ (V2 z) ∧ (V1 z) × (V2 z) ⊆ U
  by(rule bchoice)
  then obtain V2 where hv12:∧z. z∈U ⇒ openin (topology-generated-by O)
  (V1 z) ∧ openin (topology-generated-by U) (V2 z) ∧ fst z ∈ (V1 z) ∧ snd z ∈ (V2
  z) ∧ (V1 z) × (V2 z) ⊆ U
  by auto
  hence 1:U = (∪z∈U. (V1 z) × (V2 z))
  by auto
  have openin (topology-generated-by {U × V | U V. U ∈ O ∧ V ∈ U}) (∪z∈U.
  (V1 z) × (V2 z))
  proof(rule openin-Union)
    show ∧S. S ∈ (λz. V1 z × V2 z) ‘ U ⇒ openin (topology-generated-by {U
  × V | U V. U ∈ O ∧ V ∈ U}) S
    proof safe
      fix x y
      assume h:(x,y) ∈ U
      then have generate-topology-on O (V1 (x,y))
      using hv12 by(auto simp: openin-topology-generated-by-iff)
      thus openin (topology-generated-by {U × V | U V. U ∈ O ∧ V ∈ U}) (V1
  (x, y) × V2 (x, y))
    proof induction
      case Empty
      then show ?case by auto
    next
      case (Int a b)
      thus ?case

```

```

      by (auto simp: Sigma-Int-distrib1)
    next
      case (UN K)
      then have openin (topology-generated-by {U × V | U V. U ∈ O ∧ V ∈ U})
        (⋃ {k × V2 (x, y) | k. k ∈ K})
        by auto
      moreover have (⋃ {k × V2 (x, y) | k. k ∈ K}) = (⋃ K × V2 (x, y))
        by blast
      ultimately show ?case by simp
    next
      case ho:(Basis s)
      have generate-topology-on U (V2 (x,y))
        using h hv12 by(auto simp: openin-topology-generated-by-iff)
      thus ?case
      proof induction
        case Empty
        then show ?case by auto
      next
        case (Int a b)
        then show ?case
          by (auto simp: Sigma-Int-distrib2)
      next
        case (UN K)
        then have openin (topology-generated-by {U × V | U V. U ∈ O ∧ V ∈
U}) (⋃ {s × k | k. k ∈ K})
          by auto
        moreover have (⋃ {s × k | k. k ∈ K}) = s × ⋃ K
          by blast
        ultimately show ?case by simp
      next
        case (Basis s')
        then show ?case
          using ho by(auto intro!: topology-generated-by-Basis)
      qed
    qed
  qed
  thus openin (topology-generated-by {U × V | U V. U ∈ O ∧ V ∈ U}) U
    using 1 by auto
qed

```

lemma *prod-topology-generated-by-open:*
prod-topology S S' = topology-generated-by {U × V | U V. openin S U ∧ openin S' V}
 using *prod-topology-generated-by[of {U | U. openin S U} {U | U. openin S' U}]*
topology-generated-by-open[of S, symmetric] topology-generated-by-open[of S']
 by *auto*

lemma *product-topology-cong:*

assumes $\bigwedge i. i \in I \implies S i = K i$
shows *product-topology* $S I = \text{product-topology } K I$
proof –
have $1 : \{\prod_E i \in I. X i \mid X. (\forall i. \text{openin } (S i) (X i)) \wedge \text{finite } \{i. X i \neq \text{topspace } (S i)\}\} \subseteq \{\prod_E i \in I. X i \mid X. (\forall i. \text{openin } (K i) (X i)) \wedge \text{finite } \{i. X i \neq \text{topspace } (K i)\}\}$ **if** $\bigwedge i. i \in I \implies S i = K i$ **for** $S K :: - \implies 'b \text{ topology}$
proof
fix x
assume $hx : x \in \{\prod_E i \in I. X i \mid X. (\forall i. \text{openin } (S i) (X i)) \wedge \text{finite } \{i. X i \neq \text{topspace } (S i)\}\}$
then obtain X **where** hX :
 $x = (\prod_E i \in I. X i) \wedge i. \text{openin } (S i) (X i) \text{ finite } \{i. X i \neq \text{topspace } (S i)\}$
by *auto*
define X' **where** $X' \equiv (\lambda i. \text{if } i \in I \text{ then } X i \text{ else } \text{topspace } (K i))$
have $x = (\prod_E i \in I. X' i)$
by (*auto simp: hX(1) X'-def PiE-def Pi-def*)
moreover have $\text{finite } \{i. X' i \neq \text{topspace } (K i)\}$
using *that* **by** (*auto intro!: finite-subset[OF - hX(3)] simp: X'-def*)
moreover have $\text{openin } (K i) (X' i)$ **for** i
using $hX(2)[\text{of } i]$ *that* $[\text{of } i]$ **by** (*auto simp: X'-def*)
ultimately show $x \in \{\prod_E i \in I. X i \mid X. (\forall i. \text{openin } (K i) (X i)) \wedge \text{finite } \{i. X i \neq \text{topspace } (K i)\}\}$
by (*auto intro!: exI[where x=X]*)
qed
have $\{\prod_E i \in I. X i \mid X. (\forall i. \text{openin } (S i) (X i)) \wedge \text{finite } \{i. X i \neq \text{topspace } (S i)\}\} = \{\prod_E i \in I. X i \mid X. (\forall i. \text{openin } (K i) (X i)) \wedge \text{finite } \{i. X i \neq \text{topspace } (K i)\}\}$
using $1[\text{of } S K] 1[\text{of } K S]$ *assms* **by** *auto*
thus *?thesis*
by (*simp add: product-topology-def*)
qed

lemma *topology-generated-by-without-empty*:
topology-generated-by $\mathcal{O} = \text{topology-generated-by } \{ U \in \mathcal{O}. U \neq \{\} \}$

proof (*rule topology-generated-by-eq*)
fix U
show $U \in \mathcal{O} \implies \text{openin } (\text{topology-generated-by } \{ U \in \mathcal{O}. U \neq \{\} \}) U$
by (*cases U = \{\}*) (*simp-all add: topology-generated-by-Basis*)
qed (*simp add: topology-generated-by-Basis*)

lemma *topology-from-bij*:

assumes *bij-betw* $f A (\text{topspace } S)$
shows *homeomorphic-map* (*pullback-topology* $A f S$) $S f \text{topspace}$ (*pullback-topology* $A f S$) = A

proof –

note $h = \text{bij-betw-imp-surj-on}[OF \text{ assms}] \text{bij-betw-inv-into-left}[OF \text{ assms}] \text{bij-betw-inv-into-right}[OF \text{ assms}]$

then show $[\text{simp}]: \text{topspace } (\text{pullback-topology } A f S) = A$
by (*auto simp: topology-pullback-topology*)

show *homeomorphic-map (pullback-topology A f S) S f*
by(*auto simp: homeomorphic-map-maps homeomorphic-maps-def h continuous-map-pullback[OF continuous-map-id,simplified] inv-into-into intro!: exI[where x=inv-into A f] continuous-map-pullback'[where f=f]*) (*metis (mono-tags, opaque-lifting) comp-apply continuous-map-eq continuous-map-id h(3) id-apply*)
qed

lemma *openin-pullback-topology'*:
assumes *bij-betw f A (topspace S)*
shows *openin (pullback-topology A f S) u \longleftrightarrow (openin S (f ' u)) \wedge u \subseteq A*
unfolding *openin-pullback-topology*
proof *safe*
fix *U*
assume *h:openin S U u = f -' U \cap A*
from *openin-subset[OF this(1)] assms*
have [*simp*]:*f ' (f -' U \cap A) = U*
by(*auto simp: image-def vimage-def bij-betw-def*)
show *openin S (f ' (f -' U \cap A))*
by(*simp add: h*)
next
assume *openin S (f ' u) u \subseteq A*
with *assms* **show** $\exists U. \text{openin } S \ U \wedge u = f -' U \cap A$
by(*auto intro!: exI[where x=f ' u] simp: bij-betw-def inj-on-def*)
qed

1.1.2 Isolated Point

definition *isolated-points-of :: 'a topology \Rightarrow 'a set \Rightarrow 'a set (infixr isolated'-points'-of 80) where*

X isolated-points-of A \equiv {x \in topspace X \cap A. x \notin X derived-set-of A}

lemma *isolated-points-of-eq:*

X isolated-points-of A = {x \in topspace X \cap A. $\exists U. x \in U \wedge \text{openin } X \ U \wedge U \cap (A - \{x\}) = \{\}$ }

unfolding *isolated-points-of-def* **by**(*auto simp: in-derived-set-of*)

lemma *in-isolated-points-of:*

x \in X isolated-points-of A \longleftrightarrow x \in topspace X \wedge x \in A \wedge ($\exists U. x \in U \wedge \text{openin } X \ U \wedge U \cap (A - \{x\}) = \{\}$)

by(*simp add: isolated-points-of-eq*)

lemma *derived-set-of-eq:*

x \in X derived-set-of A \longleftrightarrow x \in X closure-of (A - {x})

by(*auto simp: in-derived-set-of in-closure-of*)

1.1.3 Perfect Set

definition *perfect-set :: 'a topology \Rightarrow 'a set \Rightarrow bool where*

perfect-set X A \longleftrightarrow closedin X A \wedge X isolated-points-of A = $\{\}$

abbreviation *perfect-space* $X \equiv \text{perfect-set } X \text{ (topspace } X)$

lemma *perfect-space-euclidean*: *perfect-space* (*euclidean* :: 'a :: *perfect-space topology*)
by(*auto simp: isolated-points-of-def perfect-set-def derived-set-of-eq closure-interior*)

lemma *perfect-setI*:
assumes *closedin* $X A$
and $\bigwedge x T. \llbracket x \in A; x \in T; \text{openin } X T \rrbracket \implies \exists y \neq x. y \in T \wedge y \in A$
shows *perfect-set* $X A$
using *assms* **by**(*simp add: perfect-set-def isolated-points-of-def in-derived-set-of*)
blast

lemma *perfect-spaceI*:
assumes $\bigwedge x T. \llbracket x \in T; \text{openin } X T \rrbracket \implies \exists y \neq x. y \in T$
shows *perfect-space* X
using *assms* **by**(*auto intro!: perfect-setI*) (*meson in-mono openin-subset*)

lemma *perfect-setD*:
assumes *perfect-set* $X A$
shows *closedin* $X A A \subseteq \text{topspace } X \bigwedge x T. \llbracket x \in A; x \in T; \text{openin } X T \rrbracket \implies$
 $\exists y \neq x. y \in T \wedge y \in A$
using *assms* *closedin-subset[of X A]* **by**(*simp-all add: perfect-set-def isolated-points-of-def in-derived-set-of*) *blast*

lemma *perfect-space-perfect*:
perfect-set euclidean (*UNIV* :: 'a :: *perfect-space set*)
by(*auto simp: perfect-set-def in-isolated-points-of*) (*metis Int-Diff inf-top.right-neutral insert-Diff not-open-singleton*)

lemma *perfect-set-subtopology*:
assumes *perfect-set* $X A$
shows *perfect-space* (*subtopology* $X A$)
using *perfect-setD[OF assms]* **by**(*auto intro!: perfect-setI simp: inf.absorb-iff2 openin-subtopology*)

1.1.4 Bases and Sub-Bases in Abstract Topology

definition *subbase-in* :: [*'a topology, 'a set set*] \Rightarrow *bool* **where**
subbase-in $S \mathcal{O} \iff S = \text{topology-generated-by } \mathcal{O}$

definition *base-in* :: [*'a topology, 'a set set*] \Rightarrow *bool* **where**
base-in $S \mathcal{O} \iff (\forall U. \text{openin } S U \iff (\exists \mathcal{U}. U = \bigcup \mathcal{U} \wedge \mathcal{U} \subseteq \mathcal{O}))$

lemma *second-countable-base-in*: *second-countable* $S \iff (\exists \mathcal{O}. \text{countable } \mathcal{O} \wedge \text{base-in } S \mathcal{O})$

proof –

have [*simp*]: $\bigwedge \mathcal{B}. (\text{openin } S = \text{arbitrary union-of } (\lambda x. x \in \mathcal{B})) \iff (\forall U. \text{openin } S U \iff (\exists \mathcal{U}. U = \bigcup \mathcal{U} \wedge \mathcal{U} \subseteq \mathcal{B}))$

by(simp add: arbitrary-def union-of-def fun-eq-iff) metis
show ?thesis
by(auto simp: second-countable base-in-def)
qed

definition zero-dimensional :: 'a topology \Rightarrow bool **where**
zero-dimensional $S \iff (\exists \mathcal{O}. \text{base-in } S \ \mathcal{O} \wedge (\forall u \in \mathcal{O}. \text{openin } S \ u \wedge \text{closedin } S \ u))$

lemma openin-base:
assumes base-in $S \ \mathcal{O} \ U = \bigcup \mathcal{U}$ **and** $\mathcal{U} \subseteq \mathcal{O}$
shows openin $S \ U$
using assms **by**(auto simp: base-in-def)

lemma base-is-subbase:
assumes base-in $S \ \mathcal{O}$
shows subbase-in $S \ \mathcal{O}$
unfolding subbase-in-def topology-eq openin-topology-generated-by-iff
proof safe
fix U
assume openin $S \ U$
then obtain \mathcal{U} **where** $hu: U = \bigcup \mathcal{U} \ \mathcal{U} \subseteq \mathcal{O}$
using assms **by**(auto simp: base-in-def)
thus generate-topology-on $\mathcal{O} \ U$
by(auto intro!: generate-topology-on.UN) (auto intro!: generate-topology-on.Basis)
next
fix U
assume generate-topology-on $\mathcal{O} \ U$
then show openin $S \ U$
proof induction
case (Basis s)
then show ?case
using openin-base[OF assms, of $s \ \{s\}$]
by auto
qed auto
qed

lemma subbase-in-subset:
assumes subbase-in $S \ \mathcal{O}$ **and** $U \in \mathcal{O}$
shows $U \subseteq \text{topspace } S$
using assms(1)[simplified subbase-in-def] topology-generated-by-topspace assms
by auto

lemma subbase-in-openin:
assumes subbase-in $S \ \mathcal{O}$ **and** $U \in \mathcal{O}$
shows openin $S \ U$
using assms **by**(simp add: subbase-in-def openin-topology-generated-by-iff generate-topology-on.Basis)

lemma base-in-subset:

assumes *base-in* $S \ \mathcal{O}$ **and** $U \in \mathcal{O}$
shows $U \subseteq \text{topspace } S$
using *subbase-in-subset*[*OF base-is-subbase*[*OF assms*(1)] *assms*(2)] .

lemma *base-in-openin*:
assumes *base-in* $S \ \mathcal{O}$ **and** $U \in \mathcal{O}$
shows *openin* $S \ U$
using *subbase-in-openin*[*OF base-is-subbase*[*OF assms*(1)] *assms*(2)] .

lemma *base-in-def2*:
assumes $\bigwedge U. U \in \mathcal{O} \implies \text{openin } S \ U$
shows *base-in* $S \ \mathcal{O} \iff (\forall U. \text{openin } S \ U \longrightarrow (\forall x \in U. \exists W \in \mathcal{O}. x \in W \wedge W \subseteq U))$

proof

assume $h:\text{base-in } S \ \mathcal{O}$
show $\forall U. \text{openin } S \ U \longrightarrow (\forall x \in U. \exists W \in \mathcal{O}. x \in W \wedge W \subseteq U)$
proof safe
fix $U \ x$
assume $h':\text{openin } S \ U \ x \in U$
then obtain \mathcal{U} **where** $hu: U = \bigcup \mathcal{U} \ \mathcal{U} \subseteq \mathcal{O}$
using h **by** (*auto simp: base-in-def*)
then obtain W **where** $x \in W \ W \in \mathcal{U}$
using $h'(\mathcal{Q})$ **by** *blast*
thus $\exists W \in \mathcal{O}. x \in W \wedge W \subseteq U$
using hu **by** (*auto intro!: bexI[where x=W]*)

qed

next

assume $h:\forall U. \text{openin } S \ U \longrightarrow (\forall x \in U. \exists W \in \mathcal{O}. x \in W \wedge W \subseteq U)$
show *base-in* $S \ \mathcal{O}$
unfolding *base-in-def*
proof safe
fix U
assume *openin* $S \ U$
then have $\forall x \in U. \exists W. W \in \mathcal{O} \wedge x \in W \wedge W \subseteq U$
using h **by** *blast*
hence $\exists W. \forall x \in U. W \ x \in \mathcal{O} \wedge x \in W \ x \wedge W \ x \subseteq U$
by (*rule bchoice*)
then obtain W **where** $hw:$
 $\forall x \in U. W \ x \in \mathcal{O} \wedge x \in W \ x \wedge W \ x \subseteq U$ **by** *auto*
thus $\exists \mathcal{U}. U = \bigcup \mathcal{U} \ \mathcal{U} \subseteq \mathcal{O}$
by (*auto intro!: exI[where x=W ' U]*)

next

fix $U \ \mathcal{U}$
show $\mathcal{U} \subseteq \mathcal{O} \implies \text{openin } S \ (\bigcup \mathcal{U})$
using *assms* **by** *auto*

qed

qed

lemma *base-in-def2'*:

$base\text{-}in\ S\ \mathcal{O} \iff (\forall b \in \mathcal{O}. openin\ S\ b) \wedge (\forall x. openin\ S\ x \implies (\exists B' \subseteq \mathcal{O}. \bigcup B' = x))$

proof

assume $h: base\text{-}in\ S\ \mathcal{O}$

show $(\forall b \in \mathcal{O}. openin\ S\ b) \wedge (\forall x. openin\ S\ x \implies (\exists B' \subseteq \mathcal{O}. \bigcup B' = x))$

proof (*rule conjI*)

show $\forall b \in \mathcal{O}. openin\ S\ b$

using $openin\text{-}base[OF\ h, of\ \{-\}]$ **by** *auto*

next

show $\forall x. openin\ S\ x \implies (\exists B' \subseteq \mathcal{O}. \bigcup B' = x)$

using h **by** (*auto simp: base-in-def*)

qed

next

assume $h: (\forall b \in \mathcal{O}. openin\ S\ b) \wedge (\forall x. openin\ S\ x \implies (\exists B' \subseteq \mathcal{O}. \bigcup B' = x))$

show $base\text{-}in\ S\ \mathcal{O}$

unfolding $base\text{-}in\text{-}def$

proof *safe*

fix U

assume $openin\ S\ U$

then obtain B' **where** $B' \subseteq \mathcal{O} \bigcup B' = U$

using h **by** *blast*

thus $\exists \mathcal{U}. U = \bigcup \mathcal{U} \wedge \mathcal{U} \subseteq \mathcal{O}$

by (*auto intro!: exI[where $x=B'$]*)

next

fix $U\ \mathcal{U}$

show $\mathcal{U} \subseteq \mathcal{O} \implies openin\ S\ (\bigcup \mathcal{U})$

using h **by** *auto*

qed

qed

corollary $base\text{-}in\text{-}in\text{-}subset$:

assumes $base\text{-}in\ S\ \mathcal{O}\ openin\ S\ u\ x \in u$

shows $\exists v \in \mathcal{O}. x \in v \wedge v \subseteq u$

using $assms\ base\text{-}in\text{-}def2\ base\text{-}in\text{-}def2'$ **by** *fastforce*

lemma $base\text{-}in\text{-}without\text{-}empty$:

assumes $base\text{-}in\ S\ \mathcal{O}$

shows $base\text{-}in\ S\ \{U \in \mathcal{O}. U \neq \{\}\}$

unfolding $base\text{-}in\text{-}def2'$

proof *safe*

fix x

assume $x \in \mathcal{O} \neg openin\ S\ x$

thus $\bigwedge y. y \in \{\}$

using $base\text{-}in\text{-}openin[OF\ assms\ \langle x \in \mathcal{O} \rangle]$ **by** *simp*

next

fix x

assume $openin\ S\ x$

then obtain B' **where** $B' \subseteq \mathcal{O} \bigcup B' = x$

using $assms$ **by** (*simp add: base-in-def2'*) *metis*

thus $\exists B' \subseteq \{U \in \mathcal{O}. U \neq \{\}\}. \bigcup B' = x$
 by(auto intro!: exI[where x={y ∈ B'. y ≠ {}}])
 qed

lemma *second-countable-ex-without-empty*:
 assumes *second-countable S*
 shows $\exists \mathcal{O}. \text{countable } \mathcal{O} \wedge \text{base-in } S \ \mathcal{O} \wedge (\forall U \in \mathcal{O}. U \neq \{\})$
proof –
 obtain \mathcal{O} where *countable* \mathcal{O} *base-in* $S \ \mathcal{O}$
 using *assms second-countable-base-in* by *blast*
 thus ?thesis
 by(auto intro!: exI[where x={U ∈ O. U ≠ {}}] *base-in-without-empty*)
 qed

lemma *subtopology-subbase-in*:
 assumes *subbase-in S O*
 shows *subbase-in* (*subtopology S T*) $\{T \cap U \mid U. U \in \mathcal{O}\}$
 using *assms subtopology-generated-by*
 by(auto simp: *subbase-in-def*)

lemma *subtopology-base-in*:
 assumes *base-in S O*
 shows *base-in* (*subtopology S T*) $\{T \cap U \mid U. U \in \mathcal{O}\}$
 unfolding *base-in-def*
proof
 fix L
 show *openin* (*subtopology S T*) $L = (\exists \mathcal{U}. L = \bigcup \mathcal{U} \wedge \mathcal{U} \subseteq \{T \cap U \mid U. U \in \mathcal{O}\})$

proof
 assume *openin* (*subtopology S T*) L
 then obtain T' where *ht*:
 openin $S \ T' \ L = T' \cap T$
 by(auto simp: *openin-subtopology*)
 then obtain \mathcal{U} where *hu*:
 $T' = (\bigcup \mathcal{U}) \cap T$
 using *assms* by(auto simp: *base-in-def*)
 show $\exists \mathcal{U}. L = \bigcup \mathcal{U} \wedge \mathcal{U} \subseteq \{T \cap U \mid U. U \in \mathcal{O}\}$
 using *hu ht* by(auto intro!: exI[where x={T ∩ U | U. U ∈ U}])
next
 assume $\exists \mathcal{U}. L = \bigcup \mathcal{U} \wedge \mathcal{U} \subseteq \{T \cap U \mid U. U \in \mathcal{O}\}$
 then obtain \mathcal{U} where *hu*: $L = \bigcup \mathcal{U} \cap T$
 by *auto*
 hence $\forall U \in \mathcal{U}. \exists U' \in \mathcal{O}. U = T \cap U'$ by *blast*
 then obtain k where *hk*: $\bigwedge U. U \in \mathcal{U} \implies k \ U \in \mathcal{O} \wedge U. U \in \mathcal{U} \implies U = T \cap k \ U$
 by *metis*
 hence $L = \bigcup \{T \cap k \ U \mid U. U \in \mathcal{U}\}$
 using *hu* by *auto*
 also have $\dots = \bigcup \{k \ U \mid U. U \in \mathcal{U}\} \cap T$ by *auto*

finally have $1:L = \bigcup \{k U \mid U. U \in \mathcal{U}\} \cap T$.
moreover have $openin S (\bigcup \{k U \mid U. U \in \mathcal{U}\})$
using $hu hk assms$ **by**($auto simp: base-in-def$)
ultimately show $openin (subtopology S T) L$
by($auto intro!: exI[\text{where } x=\bigcup \{k U \mid U. U \in \mathcal{U}\}] simp: openin-subtopology$)
qed
qed

lemma *second-countable-subtopology*:
assumes *second-countable S*
shows *second-countable (subtopology S T)*
proof –
obtain \mathcal{O} **where** *countable O base-in S O*
using *assms second-countable-base-in* **by** *blast*
thus *?thesis*
by($auto intro!: exI[\text{where } x=\{T \cap U \mid U. U \in \mathcal{O}\}] simp: second-countable-base-in$
Setcompr-eq-image dest: subtopology-base-in)
qed

lemma *open-map-with-base*:
assumes *base-in S O* $\wedge A. A \in \mathcal{O} \implies openin S' (f ' A)$
shows *open-map S S' f*
unfolding *open-map-def*
proof *safe*
fix U
assume $openin S U$
then obtain \mathcal{U} **where** $U = \bigcup \mathcal{U} \mathcal{U} \subseteq \mathcal{O}$
using *assms(1)* **by**($auto simp: base-in-def$)
hence $f ' U = \bigcup \{f ' A \mid A. A \in \mathcal{U}\}$ **by** *blast*
also have $openin S' \dots$
using *assms(2)* $\langle \mathcal{U} \subseteq \mathcal{O} \rangle$ **by** *auto*
finally show $openin S' (f ' U)$.
qed

Construct a base from a subbase.

lemma *finite'-intersection-of-idempot* [*simp*]:
 $finite' \text{ intersection-of } finite' \text{ intersection-of } P = finite' \text{ intersection-of } P$
proof
fix A
show $(finite' \text{ intersection-of } finite' \text{ intersection-of } P) A = (finite' \text{ intersection-of } P) A$
proof
assume $(finite' \text{ intersection-of } finite' \text{ intersection-of } P) A$
then obtain \mathcal{U} **where** $\mathcal{U}: finite' \mathcal{U} \wedge \mathcal{U} \subseteq Collect (finite' \text{ intersection-of } P) \wedge$
 $\bigcap \mathcal{U} = A$
by($auto simp: intersection-of-def$)
hence $\forall U \in \mathcal{U}. \exists U'. finite' U' \wedge U' \subseteq Collect P \wedge \bigcap U' = U$
by($auto simp: intersection-of-def$)
then obtain \mathcal{U}' **where** \mathcal{U}' :

$\bigwedge U. U \in \mathcal{U} \implies \text{finite}' (U' U) \bigwedge U. U \in \mathcal{U} \implies U' U \subseteq \text{Collect } P \bigwedge U. U \in \mathcal{U} \implies \bigcap (U' U) = U$
 by *metis*
 have 1: $\bigcap (\bigcup (U' U)) = A$
 using $U' U(3)$ by *blast*
 show (*finite'* *intersection-of* P) A
 unfolding *intersection-of-def*
 using $U' U(1,2)$ 1 by(*auto intro!*: $\text{exI}[\text{where } x = \bigcup U \in \mathcal{U}. U' U]$)
 qed(*rule finite'-intersection-of-inc*)
 qed

lemma *finite'-intersection-of-countable*:

assumes *countable* \mathcal{O}
 shows *countable* ($\text{Collect} (\text{finite}' \text{ intersection-of } (\lambda x. x \in \mathcal{O}))$)
 proof –
 have $\text{Collect} (\text{finite}' \text{ intersection-of } (\lambda x. x \in \mathcal{O})) = (\bigcup i \in \{\mathcal{O}'. \mathcal{O}' \neq \{\}\} \wedge \text{finite } \mathcal{O}' \wedge \mathcal{O}' \subseteq \mathcal{O}). \{\bigcap i\}$
 by(*auto simp: intersection-of-def*)
 also have *countable* ...
 using *countable-Collect-finite-subset*[*OF assms*]
 by(*auto intro!*: *countable-UN*[of $\{\mathcal{O}'. \mathcal{O}' \neq \{\}\} \wedge \text{finite } \mathcal{O}' \wedge \mathcal{O}' \subseteq \mathcal{O}\} \lambda \mathcal{O}'. \{\bigcap \mathcal{O}'\}$])
 (*auto intro!*: *countable-subset*[of $\{\mathcal{O}'. \mathcal{O}' \neq \{\}\} \wedge \text{finite } \mathcal{O}' \wedge \mathcal{O}' \subseteq \mathcal{O}\} \{A. \text{finite } A \wedge A \subseteq \mathcal{O}\}$])
 finally show ?*thesis* .
 qed

lemma *finite'-intersection-of-openin*:

assumes (*finite'* *intersection-of* $(\lambda x. x \in \mathcal{O})$) U
 shows *openin* (*topology-generated-by* \mathcal{O}) U
 unfolding *openin-topology-generated-by-iff*
 using *assms* by(*auto simp: generate-topology-on-eq arbitrary-union-of-inc*)

lemma *topology-generated-by-finite-intersections*:

topology-generated-by $\mathcal{O} = \text{topology-generated-by} (\text{Collect} (\text{finite}' \text{ intersection-of } (\lambda x. x \in \mathcal{O})))$
 unfolding *topology-eq openin-topology-generated-by-iff* by(*simp add: generate-topology-on-eq*)

lemma *base-from-subbase*:

assumes *subbase-in* S \mathcal{O}
 shows *base-in* S ($\text{Collect} (\text{finite}' \text{ intersection-of } (\lambda x. x \in \mathcal{O}))$)
 unfolding *subbase-in-def base-in-def* *assms*[*simplified subbase-in-def*] *openin-topology-generated-by-iff*
 by(*auto simp: arbitrary-def union-of-def generate-topology-on-eq*)

lemma *countable-base-from-countable-subbase*:

assumes *countable* \mathcal{O} and *subbase-in* S \mathcal{O}
 shows *second-countable* S
 using *finite'-intersection-of-countable*[*OF assms(1)*] *base-from-subbase*[*OF assms(2)*]
 by(*auto simp: second-countable-base-in*)

lemma *prod-topology-second-countable*:
assumes *second-countable S and second-countable S'*
shows *second-countable (prod-topology S S')*
proof –
obtain $\mathcal{O} \ \mathcal{O}'$ **where** *ho*:
countable O base-in S O countable O' base-in S' O'
using *assms by(auto simp: second-countable-base-in)*
show *?thesis*
proof(*rule countable-base-from-countable-subbase[where O={ U × V | U V. U*
 $\in \mathcal{O} \wedge V \in \mathcal{O}'$])
have $\{U \times V \mid U V. U \in \mathcal{O} \wedge V \in \mathcal{O}'\} = (\lambda(U,V). U \times V) \text{ ` } (\mathcal{O} \times \mathcal{O}')$
by *auto*
also have *countable ...*
using *ho(1,3) by auto*
finally show *countable {U × V | U V. U ∈ O ∧ V ∈ O'}* .
next
show *subbase-in (prod-topology S S') {U × V | U V. U ∈ O ∧ V ∈ O'}*
using *base-is-subbase[OF ho(2)] base-is-subbase[OF ho(4)]*
by(*simp add: subbase-in-def prod-topology-generated-by*)
qed
qed

Abstract version of the theorem $\exists K. \textit{topological-basis } K \wedge \textit{countable } K \wedge$
 $(\forall k \in K. \exists X. k = \Pi_{i \in I} UNIV X \wedge (\forall i. \textit{open } (X i)) \wedge \textit{finite } \{i. X i \neq$
 $UNIV\})$.

lemma *product-topology-countable-base-in*:
assumes *countable I and $\bigwedge i. i \in I \implies \textit{second-countable } (S i)$*
shows $\exists \mathcal{O}'. \textit{countable } \mathcal{O}' \wedge \textit{base-in } (\textit{product-topology } S I) \ \mathcal{O}' \wedge$
 $(\forall k \in \mathcal{O}'. \exists X. k = (\Pi_{i \in I} X i) \wedge (\forall i. \textit{openin } (S i) (X i)) \wedge \textit{finite}$
 $\{i. X i \neq \textit{topspace } (S i)\} \wedge \{i. X i \neq \textit{topspace } (S i)\} \subseteq I)$
proof –
obtain \mathcal{O} **where** *ho*:
 $\bigwedge i. i \in I \implies \textit{countable } (\mathcal{O} i) \wedge i. i \in I \implies \textit{base-in } (S i) (\mathcal{O} i)$
using *assms(2)[simplified second-countable-base-in] bymetis*
show *?thesis*
unfolding *second-countable-base-in*
proof(*intro exI[where x={ $\Pi_{i \in I} U i \mid U. \textit{finite } \{i \in I. U i \neq \textit{topspace } (S i)\}$*
 $\wedge (\forall i \in \{i \in I. U i \neq \textit{topspace } (S i)\}. U i \in \mathcal{O} i)\}$ *conjI*)
show *countable { $\Pi_{i \in I} U i \mid U. \textit{finite } \{i \in I. U i \neq \textit{topspace } (S i)\} \wedge (\forall i \in \{i \in I.$*
 $U i \neq \textit{topspace } (S i)\}. U i \in \mathcal{O} i)\}$
(is countable ?X)
proof –
have $?X = \{\Pi_{i \in I} U i \mid U. \textit{finite } \{i \in I. U i \neq \textit{topspace } (S i)\} \wedge (\forall i \in \{i \in I.$
 $U i \neq \textit{topspace } (S i)\}. U i \in \mathcal{O} i) \wedge (\forall i \in (UNIV - I). U i = \{\textit{undefined}\})\}$
(is - = ?Y)
proof (*rule set-eqI*)
show $\bigwedge x. x \in ?X \longleftrightarrow x \in ?Y$
proof

```

fix  $x$ 
assume  $x \in ?X$ 
then obtain  $U$  where  $hu$ :
 $x = (\prod_E i \in I. U i)$  finite  $\{i \in I. U i \neq \text{topspace } (S i)\}$   $(\forall i \in \{i \in I. U i \neq$ 
topspace  $(S i)\}. U i \in \mathcal{O} i)$ 
  by auto
define  $U'$  where  $U' i \equiv$  (if  $i \in I$  then  $U i$  else  $\{\text{undefined}\})$  for  $i$ 
have  $x = (\prod_E i \in I. U' i)$ 
  using  $hu(1)$  by(auto simp: U'-def PiE-def extensional-def Pi-def)
  moreover have finite  $\{i \in I. U' i \neq \text{topspace } (S i)\}$   $(\forall i \in \{i \in I. U' i \neq$ 
topspace  $(S i)\}. U' i \in \mathcal{O} i) \forall i \in (UNIV - I). U' i = \{\text{undefined}\}$ 
  using  $hu(2,3)$  by(auto simp: U'-def) (metis (mono-tags, lifting) Collect-cong)
  ultimately show  $x \in ?Y$  by auto
qed auto
qed
also have  $\dots = (\lambda U. \prod_E i \in I. U i) \cdot \{U. \text{finite } \{i \in I. U i \neq \text{topspace } (S$ 
i)\} \wedge (\forall i \in \{i \in I. U i \neq \text{topspace } (S i)\}. U i \in \mathcal{O} i) \wedge (\forall i \in (UNIV - I). U i =
 $\{\text{undefined}\})\}$  by auto
also have countable  $\dots$ 
proof(rule countable-image)
  have  $\{U. \text{finite } \{i \in I. U i \neq \text{topspace } (S i)\} \wedge (\forall i \in \{i \in I. U i \neq \text{topspace}$ 
 $(S i)\}. U i \in \mathcal{O} i) \wedge (\forall i \in UNIV - I. U i = \{\text{undefined}\})\} = \{U. \exists I'. \text{finite } I' \wedge$ 
 $I' \subseteq I \wedge (\forall i \in I'. U i \in \mathcal{O} i) \wedge (\forall i \in (I - I'). U i = \text{topspace } (S i)) \wedge (\forall i \in UNIV$ 
 $- I. U i = \{\text{undefined}\})\}$ 
  (is  $?A = ?B$ )
  proof (rule set-eqI)
    show  $\bigwedge x. x \in ?A \longleftrightarrow x \in ?B$ 
    proof
      fix  $U$ 
      assume  $U \in \{U. \text{finite } \{i \in I. U i \neq \text{topspace } (S i)\} \wedge (\forall i \in \{i \in I. U i$ 
 $\neq \text{topspace } (S i)\}. U i \in \mathcal{O} i) \wedge (\forall i \in UNIV - I. U i = \{\text{undefined}\})\}$ 
      then show  $U \in \{U. \exists I'. \text{finite } I' \wedge I' \subseteq I \wedge (\forall i \in I'. U i \in \mathcal{O} i) \wedge$ 
 $(\forall i \in I - I'. U i = \text{topspace } (S i)) \wedge (\forall i \in UNIV - I. U i = \{\text{undefined}\})\}$ 
      by auto
    next
      fix  $U$ 
      assume assm:  $U \in \{U. \exists I'. \text{finite } I' \wedge I' \subseteq I \wedge (\forall i \in I'. U i \in \mathcal{O} i) \wedge$ 
 $(\forall i \in I - I'. U i = \text{topspace } (S i)) \wedge (\forall i \in UNIV - I. U i = \{\text{undefined}\})\}$ 
      then obtain  $I'$  where  $hi'$ :
         $\text{finite } I' \wedge I' \subseteq I \wedge \forall i \in I'. U i \in \mathcal{O} i \wedge \forall i \in I - I'. U i = \text{topspace } (S i)$ 
 $\forall i \in UNIV - I. U i = \{\text{undefined}\}$ 
        by auto
      then have  $\bigwedge i. i \in I \implies U i \neq \text{topspace } (S i) \implies i \in I'$  by auto
      hence  $\{i \in I. U i \neq \text{topspace } (S i)\} \subseteq I'$  by auto
      hence finite  $\{i \in I. U i \neq \text{topspace } (S i)\}$ 
      using  $hi'(1)$  by (simp add: rev-finite-subset)
      thus  $U \in \{U. \text{finite } \{i \in I. U i \neq \text{topspace } (S i)\} \wedge (\forall i \in \{i \in I. U i \neq$ 
 $\text{topspace } (S i)\}. U i \in \mathcal{O} i) \wedge (\forall i \in UNIV - I. U i = \{\text{undefined}\})\}$ 

```

```

    using hi' by auto
  qed
  qed
  also have ... = ( $\bigcup I' \in \{I'. \text{finite } I' \wedge I' \subseteq I\}. \{U. (\forall i \in I'. U i \in \mathcal{O} i) \wedge (\forall i \in I - I'. U i = \text{topspace } (S i)) \wedge (\forall i \in UNIV - I. U i = \{\text{undefined}\})\}$ )
    by auto
  also have countable ...
  proof(rule countable-UN[OF countable-Collect-finite-subset[OF assms(1)]])
    fix I'
    assume I'  $\in \{I'. \text{finite } I' \wedge I' \subseteq I\}$ 
    hence hi':finite I' I'  $\subseteq I$  by auto
    have  $(\lambda U i. \text{if } i \in I' \text{ then } U i \text{ else undefined}) ' \{U. (\forall i \in I'. U i \in \mathcal{O} i) \wedge (\forall i \in I - I'. U i = \text{topspace } (S i)) \wedge (\forall i \in UNIV - I. U i = \{\text{undefined}\})\} \subseteq (\Pi_{E i \in I'. \mathcal{O} i})$ 
      by auto
    moreover have countable ...
      using hi' by(auto intro!: countable-PiE ho)
    ultimately have countable  $((\lambda U i. \text{if } i \in I' \text{ then } U i \text{ else undefined}) ' \{U. (\forall i \in I'. U i \in \mathcal{O} i) \wedge (\forall i \in I - I'. U i = \text{topspace } (S i)) \wedge (\forall i \in UNIV - I. U i = \{\text{undefined}\})\})$ 
      by(simp add: countable-subset)
    moreover have inj-on  $(\lambda U i. \text{if } i \in I' \text{ then } U i \text{ else undefined}) \{U. (\forall i \in I'. U i \in \mathcal{O} i) \wedge (\forall i \in I - I'. U i = \text{topspace } (S i)) \wedge (\forall i \in UNIV - I. U i = \{\text{undefined}\})\}$ 
      (is inj-on ?f ?X)
  proof
    fix x y
    assume hxy:  $x \in ?X y \in ?X ?f x = ?f y$ 
    show  $x = y$ 
  proof
    fix i
    consider  $i \in I' \mid i \in I - I' \mid i \in UNIV - I$ 
    using hi'(2) by blast
    then show  $x i = y i$ 
  proof cases
    case i:1
    then show ?thesis
      using fun-cong[OF hxy(3),of i] by auto
  next
    case i:2
    then show ?thesis
      using hxy(1,2) by auto
  next
    case i:3
    then show ?thesis
      using hxy(1,2) by auto
  qed
  qed
  qed

```

ultimately show $\text{countable } \{U. (\forall i \in I'. U i \in \mathcal{O} i) \wedge (\forall i \in I - I'. U i = \text{topspace } (S i)) \wedge (\forall i \in \text{UNIV} - I. U i = \{\text{undefined}\})\}$
using *countable-image-inj-on* **by** *auto*
qed
finally show $\text{countable } \{U. \text{finite } \{i \in I. U i \neq \text{topspace } (S i)\} \wedge (\forall i \in \{i \in I. U i \neq \text{topspace } (S i)\}. U i \in \mathcal{O} i) \wedge (\forall i \in \text{UNIV} - I. U i = \{\text{undefined}\})\}$.
qed
finally show *?thesis*.
qed
next
show *base-in (product-topology S I) { $\prod_E i \in I. U i \mid U. \text{finite } \{i \in I. U i \neq \text{topspace } (S i)\} \wedge (\forall i \in \{i \in I. U i \neq \text{topspace } (S i)\}. U i \in \mathcal{O} i)\}$ }*
(is base-in (product-topology S I) ?X)
unfolding *base-in-def*
proof *safe*
fix *U*
assume *openin (product-topology S I) U*
then have $\forall x \in U. \exists Ux. \text{finite } \{i \in I. Ux i \neq \text{topspace } (S i)\} \wedge (\forall i \in I. \text{openin } (S i) (Ux i)) \wedge x \in \text{Pi}_E I Ux \wedge \text{Pi}_E I Ux \subseteq U$
by *(simp add: openin-product-topology-alt)*
hence $\exists Ux. \forall x \in U. \text{finite } \{i \in I. Ux x i \neq \text{topspace } (S i)\} \wedge (\forall i \in I. \text{openin } (S i) (Ux x i)) \wedge x \in \text{Pi}_E I (Ux x) \wedge \text{Pi}_E I (Ux x) \subseteq U$
by *(rule bchoice)*
then obtain *Ux where hui:*
 $\bigwedge x. x \in U \implies \text{finite } \{i \in I. Ux x i \neq \text{topspace } (S i)\} \bigwedge x i. x \in U \implies i \in I \implies \text{openin } (S i) (Ux x i) \bigwedge x. x \in U \implies x \in \text{Pi}_E I (Ux x) \bigwedge x. x \in U \implies \text{Pi}_E I (Ux x) \subseteq U$
by *fastforce*
then have $1: \forall x \in U. \forall i \in \{i \in I. Ux x i \neq \text{topspace } (S i)\}. \exists Uxj. Uxj \subseteq \mathcal{O} i \wedge Ux x i = \bigcup Uxj$
using *ho[simplified base-in-def]* **by** *(metis (no-types, lifting) mem-Collect-eq)*

have $\forall x \in U. \exists Uxj. \forall i \in \{i \in I. Ux x i \neq \text{topspace } (S i)\}. Uxj i \subseteq \mathcal{O} i \wedge Ux x i = \bigcup (Uxj i)$
by *(standard, rule bchoice) (use 1 in simp)*
hence $\exists Uxj. \forall x \in U. \forall i \in \{i \in I. Ux x i \neq \text{topspace } (S i)\}. Uxj x i \subseteq \mathcal{O} i \wedge Ux x i = \bigcup (Uxj x i)$
by *(rule bchoice)*
then obtain *Uxj where*
 $\forall x \in U. \forall i \in \{i \in I. Ux x i \neq \text{topspace } (S i)\}. Uxj x i \subseteq \mathcal{O} i \wedge Ux x i = \bigcup (Uxj x i)$
by *auto*
hence *huxj:* $\bigwedge x i. x \in U \implies i \in \{i \in I. Ux x i \neq \text{topspace } (S i)\} \implies Uxj x i \subseteq \mathcal{O} i$
 $\bigwedge x i. x \in U \implies i \in \{i \in I. Ux x i \neq \text{topspace } (S i)\} \implies Ux x i = \bigcup (Uxj x i)$
by *blast+*
show $\exists U. U = \bigcup U \wedge U \subseteq ?X$
proof *(intro exI[where x={ $\prod_E i \in I. K i \mid K. \exists x \in U. \text{finite } \{i \in I. Ux x$*

$i \neq \text{topspace } (S \ i)\} \wedge (\forall i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. K \ i \in \mathcal{U}xj \ x \ i) \wedge$
 $(\forall i \in UNIV - \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. K \ i = \text{topspace } (S \ i))\} \text{ conj}I$
show $U = \bigcup \{\prod_E \ i \in I. K \ i \mid K. \exists x \in U. \text{finite } \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\} \wedge (\forall i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. K \ i \in \mathcal{U}xj \ x \ i) \wedge (\forall i \in UNIV - \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. K \ i = \text{topspace } (S \ i))\}$
proof safe
fix x
assume $hxu: x \in U$
have $\forall i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. Ux \ x \ i = \bigcup (\mathcal{U}xj \ x \ i)$
using $hxj[OF \ hxu]$ **by blast**
hence $\forall i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. \exists Uxj. Uxj \in \mathcal{U}xj \ x \ i \wedge x \ i \in Uxj$
using $hui(3)[OF \ hxu]$ **by auto**
hence $\exists Uxj. \forall i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. Uxj \ i \in \mathcal{U}xj \ x \ i \wedge x \ i \in Uxj \ i$
by (rule bchoice)
then obtain Uxj **where** hxj' :
 $\wedge i. i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\} \implies Uxj \ i \in \mathcal{U}xj \ x \ i$
 $\wedge i. i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\} \implies x \ i \in Uxj \ i$
by auto
define K **where** $K \equiv (\lambda i. \text{if } i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\} \text{ then } Uxj \ i \text{ else } \text{topspace } (S \ i))$
have $x \in (\prod_E \ i \in I. K \ i)$
using $hxj'(2)$ $hui(3,4)[OF \ hxu]$ $\text{openin-subset}[OF \ hui(2)[OF \ hxu]]$
by (auto simp: K-def PiE-def Pi-def)
moreover have $\exists x \in U. \text{finite } \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\} \wedge (\forall i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. K \ i \in \mathcal{U}xj \ x \ i) \wedge (\forall i \in UNIV - \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. K \ i = \text{topspace } (S \ i))$
by (rule best[OF - hxu], rule conjI, simp add: hui(1)[OF \ hxu]) (use hui(2) hxu openin-subset hxj'(1) K-def in auto)
ultimately show $x \in \bigcup \{\prod_E \ i \in I. K \ i \mid K. \exists x \in U. \text{finite } \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\} \wedge (\forall i \in \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. K \ i \in \mathcal{U}xj \ x \ i) \wedge (\forall i \in UNIV - \{i \in I. Ux \ x \ i \neq \text{topspace } (S \ i)\}. K \ i = \text{topspace } (S \ i))\}$
by auto
next
fix $x \ X \ K \ u$
assume $hu: x \in (\prod_E \ i \in I. K \ i) \ u \in U \ \text{finite } \{i \in I. Ux \ u \ i \neq \text{topspace } (S \ i)\} \wedge (\forall i \in \{i \in I. Ux \ u \ i \neq \text{topspace } (S \ i)\}. K \ i \in \mathcal{U}xj \ u \ i) \wedge (\forall i \in UNIV - \{i \in I. Ux \ u \ i \neq \text{topspace } (S \ i)\}. K \ i = \text{topspace } (S \ i))$
have $\wedge i. i \in \{i \in I. Ux \ u \ i \neq \text{topspace } (S \ i)\} \implies K \ i \subseteq Ux \ u \ i$
using $hxj[OF \ hu(2)]$ $hu(4)$ **by blast**
moreover have $\wedge i. i \in I - \{i \in I. Ux \ u \ i \neq \text{topspace } (S \ i)\} \implies K \ i = Ux \ u \ i$
using $hu(5)$ **by auto**
ultimately have $\wedge i. i \in I \implies K \ i \subseteq Ux \ u \ i$
by blast
thus $x \in U$
using $hui(4)[OF \ hu(2)]$ $hu(1)$ **by blast**
qed

next
show $\{\prod_E i \in I. K i \mid K. \exists x \in U. \text{finite } \{i \in I. Ux x i \neq \text{topspace } (S i)\} \wedge$
 $(\forall i \in \{i \in I. Ux x i \neq \text{topspace } (S i)\}. K i \in \mathcal{U}xj x i) \wedge (\forall i \in UNIV - \{i \in I. Ux x$
 $i \neq \text{topspace } (S i)\}. K i = \text{topspace } (S i))\} \subseteq ?X$
proof
fix x
assume $x \in \{\prod_E i \in I. K i \mid K. \exists x \in U. \text{finite } \{i \in I. Ux x i \neq \text{topspace}$
 $(S i)\} \wedge (\forall i \in \{i \in I. Ux x i \neq \text{topspace } (S i)\}. K i \in \mathcal{U}xj x i) \wedge (\forall i \in UNIV - \{i$
 $\in I. Ux x i \neq \text{topspace } (S i)\}. K i = \text{topspace } (S i))\}$
then obtain $u K$ **where** hu :
 $x = (\prod_E i \in I. K i) \ u \in U \text{finite } \{i \in I. Ux u i \neq \text{topspace } (S i)\} \forall i \in \{i \in$
 $I. Ux u i \neq \text{topspace } (S i)\}. K i \in \mathcal{U}xj u i \forall i \in UNIV - \{i \in I. Ux u i \neq \text{topspace}$
 $(S i)\}. K i = \text{topspace } (S i)$
by *auto*
have $hksbst: \{i \in I. K i \neq \text{topspace } (S i)\} \subseteq \{i \in I. Ux u i \neq \text{topspace}$
 $(S i)\}$
using $hu(5)$ **by** *fastforce*
hence $\text{finite } \{i \in I. K i \neq \text{topspace } (S i)\}$
using $hu(3)$ **by** (*simp add: finite-subset*)
moreover have $\forall i \in \{i \in I. K i \neq \text{topspace } (S i)\}. K i \in \mathcal{O} i$
using $huxj(1)[OF hu(2)] hu(4) hksbst$
by (*meson subsetD*)
ultimately show $x \in ?X$
using $hu(1)$ **by** *auto*
qed
qed
next
fix \mathcal{U}
assume $\mathcal{U} \subseteq ?X$
have $\text{openin } (\text{product-topology } S I) u$ **if** $hu: u \in \mathcal{U}$ **for** u
proof –
have $hu': u \in ?X$
using $\langle \mathcal{U} \subseteq ?X \rangle hu$ **by** *auto*
then obtain U **where** hU :
 $u = (\prod_E i \in I. U i) \text{finite } \{i \in I. U i \neq \text{topspace } (S i)\} \forall i \in \{i \in I. U i \neq$
 $\text{topspace } (S i)\}. U i \in \mathcal{O} i$
by *auto*
define U' **where** $U' \equiv (\lambda i. \text{if } i \in \{i \in I. U i \neq \text{topspace } (S i)\} \text{ then } U i$
 $\text{else } \text{topspace } (S i))$
have $hU': u = (\prod_E i \in I. U' i)$
by(*auto simp: hU(1) U'-def PiE-def Pi-def*)
have $hU\text{finite} : \text{finite } \{i. U' i \neq \text{topspace } (S i)\}$
using $hU(2)$ **by**(*auto simp: U'-def*)
have $hUoi: \forall i \in \{i. U' i \neq \text{topspace } (S i)\}. U' i \in \mathcal{O} i$
using $hU(3)$ **by**(*auto simp: U'-def*)
have $hUi: \forall i \in \{i. U' i \neq \text{topspace } (S i)\}. i \in I$
using $hU(2)$ **by**(*auto simp: U'-def*)
have $\text{hallopen: openin } (S i) (U' i)$ **for** i
proof –

```

    consider  $i \in \{i. U' i \neq \text{topspace } (S i)\} \mid i \notin \{i. U' i \neq \text{topspace } (S i)\}$ 
by auto
  then show ?thesis
  proof cases
    case 1
    then show ?thesis
      using  $hUoi \text{ ho}(2)[\text{of } i] \text{ base-in-openin}[\text{of } S i \ \mathcal{O} \ i \ U' i] \ hUi$ 
      by auto
    next
    case 2
    then have  $U' i = \text{topspace } (S i)$  by auto
    thus ?thesis by auto
  qed
  show  $\text{openin } (\text{product-topology } S \ I) \ u$ 
    using  $\text{hallopen } hU\text{finite} \text{ by}(\text{auto intro!}: \text{product-topology-basis simp}: hU')$ 
  qed
  thus  $\text{openin } (\text{product-topology } S \ I) \ (\bigcup \ \mathcal{U})$ 
  by auto
  qed
next
  show  $\forall k \in \{Pi_E \ I \ U \mid U. \text{finite } \{i \in I. U i \neq \text{topspace } (S i)\} \wedge (\forall i \in \{i \in I. U i \neq \text{topspace } (S i)\}. U i \in \mathcal{O} i)\}. \exists X. k = Pi_E \ I \ X \wedge (\forall i. \text{openin } (S i) (X i)) \wedge \text{finite } \{i. X i \neq \text{topspace } (S i)\} \wedge \{i. X i \neq \text{topspace } (S i)\} \subseteq I$ 
  proof
    fix  $k$ 
    assume  $k \in \{Pi_E \ I \ U \mid U. \text{finite } \{i \in I. U i \neq \text{topspace } (S i)\} \wedge (\forall i \in \{i \in I. U i \neq \text{topspace } (S i)\}. U i \in \mathcal{O} i)\}$ 
    then obtain  $U$  where  $hu$ :
       $k = (\Pi_E \ i \in I. U i) \text{ finite } \{i \in I. U i \neq \text{topspace } (S i)\} \forall i \in \{i \in I. U i \neq \text{topspace } (S i)\}. U i \in \mathcal{O} i$ 
    by auto
    define  $X$  where  $X \equiv (\lambda i. \text{if } i \in \{i \in I. U i \neq \text{topspace } (S i)\} \text{ then } U i \text{ else } \text{topspace } (S i))$ 
    have  $hX1: k = (\Pi_E \ i \in I. X i)$ 
      using  $hu(1) \text{ by}(\text{auto simp}: X\text{-def } PiE\text{-def } Pi\text{-def})$ 
    have  $hX2: \text{openin } (S i) (X i)$  for  $i$ 
      using  $hu(3) \text{ base-in-openin}[\text{of } S i \ - \ U i, OF \ \text{ho}(2)]$ 
      by  $(\text{auto simp}: X\text{-def})$ 
    have  $hX3: \text{finite } \{i. X i \neq \text{topspace } (S i)\}$ 
      using  $hu(2) \text{ by}(\text{auto simp}: X\text{-def})$ 
    have  $hX4: \{i. X i \neq \text{topspace } (S i)\} \subseteq I$ 
      by  $(\text{auto simp}: X\text{-def})$ 
    show  $\exists X. k = (\Pi_E \ i \in I. X i) \wedge (\forall i. \text{openin } (S i) (X i)) \wedge \text{finite } \{i. X i \neq \text{topspace } (S i)\} \wedge \{i. X i \neq \text{topspace } (S i)\} \subseteq I$ 
      using  $hX1 \ hX2 \ hX3 \ hX4 \text{ by}(\text{auto intro!}: \text{exI}[\text{where } x=X])$ 
  qed
  qed
  qed

```

lemma *product-topology-second-countable*:
assumes *countable I and $\bigwedge i. i \in I \implies \text{second-countable } (S\ i)$*
shows *second-countable (product-topology S I)*
using *product-topology-countable-base-in[OF assms(1)] assms(2)*
by(*fastforce simp: second-countable-base-in*)

lemma *second-countable-euclidean[simp]*:
second-countable (euclidean :: 'a :: second-countable-topology topology)
using *ex-countable-basis second-countable-def topological-basis-def* **by** *fastforce*

lemma *Cantor-Bendixon*:
assumes *second-countable X*
shows $\exists U\ P. \text{countable } U \wedge \text{openin } X\ U \wedge \text{perfect-set } X\ P \wedge U \cup P = \text{topspace } X \wedge U \cap P = \{\} \wedge (\forall a \neq \{\}. \text{openin } (\text{subtopology } X\ P)\ a \longrightarrow \text{uncountable } a)$
proof –
obtain \mathcal{O} **where** *o: countable \mathcal{O} base-in X \mathcal{O}*
using *assms by(auto simp: second-countable-base-in)*
define U **where** $U \equiv \bigcup \{u \in \mathcal{O}. \text{countable } u\}$
define P **where** $P \equiv \text{topspace } X - U$
have *1: countable U*
using *o(1) by(auto simp: U-def intro!: countable-UN[of - id,simplified])*
have *2: openin X U*
using *base-in-openin[OF o(2)] by(auto simp: U-def)*
have *openin-c:countable v $\longleftrightarrow v \subseteq U$ if openin X v for v*
proof
assume *countable v*
obtain \mathcal{U} **where** $v = \bigcup \mathcal{U} \ \mathcal{U} \subseteq \mathcal{O}$
using $\langle \text{openin } X\ v \rangle$ *o(2) by(auto simp: base-in-def)*
with $\langle \text{countable } v \rangle$ **have** $\bigwedge u. u \in \mathcal{U} \implies \text{countable } u$
by (*meson Sup-upper countable-subset*)
thus $v \subseteq U$
using $\langle U \subseteq \mathcal{O} \rangle$ *by(auto simp: $\langle v = \bigcup \mathcal{U} \rangle$ U-def)*
qed(*rule countable-subset[OF - 1]*)
have *3: perfect-set X P*
proof(*rule perfect-setI*)
fix $x\ T$
assume $h: x \in P\ x \in T\ \text{openin } X\ T$
have *T-unc:uncountable T*
using *openin-c[OF h(3)] h(1,2) by(auto simp: P-def)*
obtain \mathcal{U} **where** $U: T = \bigcup \mathcal{U} \ \mathcal{U} \subseteq \mathcal{O}$
using *h(3) o(2) by(auto simp: base-in-def)*
then obtain u **where** $u: u \in \mathcal{U}$ *uncountable u*
using *T-unc U-def h(3) openin-c* **by** *auto*
hence *uncountable (u - {x})* **by** *simp*
hence $\neg (u - \{x\} \subseteq U)$
using *1* **by** (*metis countable-subset*)
then obtain y **where** $y \in u - \{x\}$ $y \notin U$
by *blast*


```

thus  $\exists y. y \neq x \wedge y \in T \wedge y \in P$ 
  using  $U \text{ u base-in-subset}[OF \ o(2),of \ u]$  by(auto intro!:  $exI$ [where  $x=y$ ])
simp:P-def)
qed(use 2 P-def in auto)
have  $\not\vdash : \text{uncountable } a \text{ if } \text{openin } (\text{subtopology } X \ P) \ a \ a \neq \{\}$  for  $a$ 
proof
  assume contable:countable a
  obtain  $b$  where  $b : \text{openin } X \ b \ a = P \cap b$ 
  using  $\langle \text{openin } (\text{subtopology } X \ P) \ a \rangle$  by(auto simp: openin-subtopology)
  hence uncountable b
  using P-def openin-c that(2) by auto
  thus False
  by (metis 1 Diff-Int-distrib2 Int-absorb1 P-def b(1) b(2) contable countable-Int1
openin-subset uncountable-minus-countable)
qed
show ?thesis
  using  $1 \ 2 \ 3 \ 4$  by(auto simp: P-def)
qed

```

1.1.5 Separable Spaces

definition *dense-in* :: [*'a topology, 'a set*] \Rightarrow *bool* **where**
dense-in $S \ U \longleftrightarrow (U \subseteq \text{topspace } S \wedge (\forall V. \text{openin } S \ V \longrightarrow V \neq \{\} \longrightarrow U \cap V \neq \{\}))$

lemma *dense-in-def2*:

```

dense-in  $S \ U \longleftrightarrow (U \subseteq \text{topspace } S \wedge (S \ \text{closure-of } U) = \text{topspace } S)$ 
using dense-intersects-open by(auto simp: dense-in-def closure-of-subset-topspace
in-closure-of) auto

```

lemma *dense-in-topspace*[*simp*]: *dense-in* $S \ (\text{topspace } S)$

by(*auto simp: dense-in-def2*)

lemma *dense-in-subset*:

```

assumes dense-in  $S \ U$ 
shows  $U \subseteq \text{topspace } S$ 
using assms by(simp add: dense-in-def)

```

lemma *dense-in-nonempty*:

```

assumes  $\text{topspace } S \neq \{\}$  dense-in  $S \ U$ 
shows  $U \neq \{\}$ 
using assms by(auto simp: dense-in-def)

```

lemma *dense-inI*:

```

assumes  $U \subseteq \text{topspace } S$ 
  and  $\bigwedge V. \text{openin } S \ V \Longrightarrow V \neq \{\} \Longrightarrow U \cap V \neq \{\}$ 
shows dense-in  $S \ U$ 
using assms by(auto simp: dense-in-def)

```

lemma *dense-in-infinite*:
assumes *t1-space X infinite (topspace X) dense-in X U*
shows *infinite U*
proof
assume *fin: finite U*
then have *closedin X U*
by (*metis assms(1,3) dense-in-def t1-space-closedin-finite*)
hence *X closure-of U = U*
by (*simp add: closure-of-eq*)
thus *False*
by (*metis assms(2) assms(3) dense-in-def2 fin*)
qed

lemma *dense-in-prod*:
assumes *dense-in S U and dense-in S' U'*
shows *dense-in (prod-topology S S') (U × U')*
proof(*rule dense-inI*)
fix *V*
assume *h:openin (prod-topology S S') V V ≠ {}*
then obtain *x y where hxy:(x,y) ∈ V by auto*
then obtain *V1 V2 where hv12:*
openin S V1 openin S' V2 x ∈ V1 y ∈ V2 V1 × V2 ⊆ V
using *h(1) openin-prod-topology-alt[of S S' V] by blast*
hence *V1 ≠ {} V2 ≠ {} by auto*
hence *U ∩ V1 ≠ {} U' ∩ V2 ≠ {}*
using *assms hv12 by(auto simp: dense-in-def)*
thus *U × U' ∩ V ≠ {}*
using *hv12 by auto*
next
show *U × U' ⊆ topspace (prod-topology S S')*
using *assms by(auto simp add: dense-in-def)*
qed

lemma *separable-space-def2:separable-space S ↔ (∃ U. countable U ∧ dense-in S U)*
by(*auto simp: separable-space-def dense-in-def2*)

lemma *countable-space-separable-space*:
assumes *countable (topspace S)*
shows *separable-space S*
using *assms by(auto simp: separable-space-def2 intro!: exI[where x=topspace S])*

lemma *separable-space-prod*:
assumes *separable-space S and separable-space S'*
shows *separable-space (prod-topology S S')*
proof –
obtain *U U' where*
countable U dense-in S U countable U' dense-in S' U'

using *assms* **by**(*auto simp: separable-space-def2*)
thus *?thesis*
by(*auto intro!: exI[where x=U×U'] dense-in-prod simp: separable-space-def2*)
qed

lemma *dense-in-product:*

assumes $\bigwedge i. i \in I \implies \text{dense-in } (T i) (U i)$
shows $\text{dense-in } (\text{product-topology } T I) (\prod_{E} i \in I. U i)$
proof(*rule dense-inI*)
fix V
assume $h: \text{openin } (\text{product-topology } T I) V \ V \neq \{\}$
then obtain x **where** $hx: x \in V$ **by** *auto*
then obtain K **where** $hk:$
 $\text{finite } \{i \in I. K i \neq \text{topspace } (T i)\} \ \forall i \in I. \text{openin } (T i) (K i) \ x \in (\prod_{E} i \in I. K i)$
 $(\prod_{E} i \in I. K i) \subseteq V$
using $h(1)$ *openin-product-topology-alt[of T I V]* **by** *auto*
hence $\bigwedge i. i \in I \implies K i \neq \{\}$ **by** *auto*
hence $\bigwedge i. i \in I \implies U i \cap K i \neq \{\}$
using *assms* hk **by**(*auto simp: dense-in-def*)
hence $(\prod_{E} i \in I. U i) \cap (\prod_{E} i \in I. K i) \neq \{\}$
by (*simp add: PiE-Int PiE-eq-empty-iff*)
thus $(\prod_{E} i \in I. U i) \cap V \neq \{\}$
using hk **by** *auto*
next
show $(\prod_{E} i \in I. U i) \subseteq \text{topspace } (\text{product-topology } T I)$
using *assms* **by**(*auto simp: dense-in-def*)
qed

lemma *separable-countable-product:*

assumes *countable I* **and** $\bigwedge i. i \in I \implies \text{separable-space } (T i)$
shows $\text{separable-space } (\text{product-topology } T I)$
proof –
consider $\exists i \in I. T i = \text{trivial-topology} \mid \bigwedge i. i \in I \implies T i \neq \text{trivial-topology}$
by *auto*
thus *?thesis*
proof *cases*
case 1
then obtain i **where** $i: i \in I \ \text{topspace } (T i) = \{\}$
by *auto*
show *?thesis*
unfolding *separable-space-def2 dense-in-def*
proof(*intro exI[where x={}] conjI*)
show $\forall V. \text{openin } (\text{product-topology } T I) V \longrightarrow V \neq \{\} \longrightarrow \{\} \cap V \neq \{\}$
proof *safe*
fix $V \ x$
assume $h: \text{openin } (\text{product-topology } T I) V \ x \in V$
from i **have** $(\text{product-topology } T I) = \text{trivial-topology}$
using *product-topology-trivial-iff* **by** *auto*
with $h(1)$ **have** $V = \{\}$

```

    by simp
  thus  $x \in \{\}$ 
    using  $h(2)$  by auto
qed
qed auto
next
case 2
then have  $\exists x. \forall i \in I. x i \in \text{topspace } (T i) \exists U. \forall i \in I. \text{countable } (U i) \wedge \text{dense-in } (T i) (U i)$ 
  using  $\text{assms}(2)$  by( $\text{auto intro! bchoice simp: separable-space-def2 ex-in-conv}$ )
  then obtain  $x U$  where  $hxu$ :
     $\bigwedge i. i \in I \implies x i \in \text{topspace } (T i) \bigwedge i. i \in I \implies \text{countable } (U i) \bigwedge i. i \in I \implies \text{dense-in } (T i) (U i)$ 
  by auto
  define  $U'$  where  $U' \equiv (\lambda J i. \text{if } i \in J \text{ then } U i \text{ else } \{x i\})$ 
  show ?thesis
    unfolding  $\text{separable-space-def2}$ 
  proof(intro  $\text{exI}$  [where  $x = \bigcup \{\Pi_E i \in I. U' J i \mid J. \text{finite } J \wedge J \subseteq I\}$ ] conjI)
    have  $(\bigcup \{\Pi_E i \in I. U' J i \mid J. \text{finite } J \wedge J \subseteq I\}) = (\bigcup ((\lambda J. \Pi_E i \in I. U' J i) ' \{J. \text{finite } J \wedge J \subseteq I\}))$ 
    by auto
    also have  $\text{countable } \dots$ 
  proof(rule  $\text{countable-UN}$ )
    fix  $J$ 
    assume  $hj: J \in \{J. \text{finite } J \wedge J \subseteq I\}$ 
    have  $\text{inj-on } (\lambda f. (\lambda i \in J. f i, \lambda i \in (I-J). f i)) (\Pi_E i \in I. U' J i)$ 
    proof(rule  $\text{inj-onI}$ )
      fix  $f g$ 
      assume  $h: f \in \text{PiE } I (U' J) g \in \text{PiE } I (U' J)$ 
         $(\text{restrict } f J, \text{restrict } f (I - J)) = (\text{restrict } g J, \text{restrict } g (I - J))$ 
      then have  $\bigwedge i. i \in J \implies f i = g i \bigwedge i. i \in (I-J) \implies f i = g i$ 
        by( $\text{auto simp: restrict-def}$ ) meson+
      thus  $f = g$ 
        using  $h(1,2)$  by( $\text{auto simp: } U'\text{-def}$ ) (meson  $\text{PiE-ext}$ )
    qed
  moreover have  $\text{countable } ((\lambda f. (\lambda i \in J. f i, \lambda i \in (I-J). f i)) ' (\Pi_E i \in I. U' J i))$  (is  $\text{countable } ?K$ )
  proof -
    have  $1: ?K \subseteq (\Pi_E i \in J. U i) \times (\Pi_E i \in (I-J). \{x i\})$ 
      using  $hj$  by( $\text{auto simp: } U'\text{-def PiE-def Pi-def}$ )
    have  $2: \text{countable } \dots$ 
  proof(rule  $\text{countable-SIGMA}$ )
    show  $\text{countable } (\text{PiE } J U)$ 
      using  $hj hxu(2)$  by( $\text{auto intro! countable-PiE}$ )
  next
    have  $(\Pi_E i \in I - J. \{x i\}) = \{\lambda i \in I - J. x i\}$ 
      by( $\text{auto simp: PiE-def extensional-def restrict-def Pi-def}$ )
    thus  $\text{countable } (\Pi_E i \in I - J. \{x i\})$ 
      by  $\text{simp}$ 
  end

```

```

qed
show ?thesis
  by(rule countable-subset[OF 1 2])
qed
ultimately show countable  $(\prod_E i \in I. U' J i)$ 
  by(simp add: countable-image-inj-eq)
qed(rule countable-Collect-finite-subset[OF assms(1)])
finally show countable  $(\bigcup \{ \prod_E i \in I. U' J i \mid J. \text{finite } J \wedge J \subseteq I \})$  .
next
show dense-in (product-topology T I)  $(\bigcup \{ \prod_E i \in I. U' J i \mid J. \text{finite } J \wedge J \subseteq I \})$ 
proof(rule dense-inI)
  fix V
  assume h: openin (product-topology T I) V  $V \neq \{ \}$ 
  then obtain y where hx:  $y \in V$  by auto
  then obtain K where hk:
    finite  $\{ i \in I. K i \neq \text{topspace } (T i) \} \wedge i. i \in I \implies \text{openin } (T i) (K i) y \in$ 
 $(\prod_E i \in I. K i) (\prod_E i \in I. K i) \subseteq V$ 
    using h(1) openin-product-topology-alt[of T I V] by auto
    hence 3:  $\bigwedge i. i \in I \implies K i \neq \{ \}$  by auto
    hence 4:  $i \in \{ i \in I. K i \neq \text{topspace } (T i) \} \implies K i \cap U' \{ i \in I. K i \neq$ 
topspace (T i)  $\} i \neq \{ \}$  for i
    using hxu(3)[of i] hk(2)[of i] by(auto simp: U'-def dense-in-def)
    have  $\exists z. \forall i \in \{ i \in I. K i \neq \text{topspace } (T i) \}. z i \in K i \cap U' \{ i \in I. K i \neq$ 
topspace (T i)  $\} i$ 
    by(rule bchoice) (use 4 in auto)
    then obtain z where hz:  $\forall i \in \{ i \in I. K i \neq \text{topspace } (T i) \}. z i \in K i \cap$ 
 $U' \{ i \in I. K i \neq \text{topspace } (T i) \} i$ 
    by auto
    have 5:  $i \notin \{ i \in I. K i \neq \text{topspace } (T i) \} \implies i \in I \implies x i \in K i$  for i
    using hxu(1)[of i] by auto
    have  $(\lambda i. \text{if } i \in \{ i \in I. K i \neq \text{topspace } (T i) \} \text{ then } z i \text{ else if } i \in I \text{ then } x i$ 
else undefined  $) \in (\prod_E i \in I. U' \{ i \in I. K i \neq \text{topspace } (T i) \} i) \cap (\prod_E i \in I. K i)$ 
    using 4 5 hz by(auto simp: U'-def)
    thus  $\bigcup \{ P i_E I (U' J) \mid J. \text{finite } J \wedge J \subseteq I \} \cap V \neq \{ \}$ 
    using hk(1,4) by blast
  next
  have  $\bigwedge J. J \subseteq I \implies (\prod_E i \in I. U' J i) \subseteq \text{topspace } (\text{product-topology } T I)$ 
using hxu by(auto simp: dense-in-def U'-def PiE-def Pi-def) (metis
subsetD)
  thus  $(\bigcup \{ \prod_E i \in I. U' J i \mid J. \text{finite } J \wedge J \subseteq I \}) \subseteq \text{topspace } (\text{product-topology}$ 
T I)
  by auto
qed
qed
qed
qed

```

lemma separable-finite-product:

assumes *finite I and* $\bigwedge i. i \in I \implies \text{separable-space } (T i)$
shows *separable-space (product-topology T I)*
using *separable-countable-product[OF countable-finite[OF assms(1)]]* **assms** **by**
auto

1.1.6 G_δ Set

lemma *gdelta-inD*:
assumes *gdelta-in S A*
shows $\exists \mathcal{U}. \mathcal{U} \neq \{\}$ \wedge *countable* $\mathcal{U} \wedge (\forall b \in \mathcal{U}. \text{openin } S b) \wedge A = \bigcap \mathcal{U}$
using *assms unfolding gdelta-in-def relative-to-def intersection-of-def*
by (*metis IntD1 Int-insert-right-if1 complete-lattice-class.Inf-insert countable-insert empty-not-insert inf.absorb-iff2 mem-Collect-eq openin-topospace*)

lemma *gdelta-inD'*:
assumes *gdelta-in S A*
shows $\exists U. (\forall n::\text{nat}. \text{openin } S (U n)) \wedge A = \bigcap (\text{range } U)$
proof –
obtain \mathcal{U} **where** $h:\mathcal{U} \neq \{\}$ *countable* $\mathcal{U} \wedge b. b \in \mathcal{U} \implies \text{openin } S b \wedge A = \bigcap \mathcal{U}$
using *gdelta-inD[OF assms]* **by** *metis*
show *?thesis*
using *range-from-nat-into[OF h(1,2)] h(3,4)*
by (*auto intro!: exI[where x=from-nat-into U]*)
qed

lemma *gdelta-in-continuous-map*:
assumes *continuous-map X Y f gdelta-in Y a*
shows *gdelta-in X (f -' a \cap topspace X)*
proof –
obtain Ua **where** u :
 $Ua \neq \{\}$ *countable* $Ua \wedge b. b \in Ua \implies \text{openin } Y b \wedge a = \bigcap Ua$
using *gdelta-inD[OF assms(2)]* **by** *metis*
then have $0:f -' a \cap \text{topspace } X = \bigcap \{f -' b \cap \text{topspace } X \mid b. b \in Ua\}$
by *auto*
have $1:\{f -' b \cap \text{topspace } X \mid b. b \in Ua\} \neq \{\}$
using $u(1)$ **by** *simp*
have $2:\text{countable } \{f -' b \cap \text{topspace } X \mid b. b \in Ua\}$
using u **by** (*simp add: Setcompr-eq-image*)
have $3:\bigwedge c. c \in \{f -' b \cap \text{topspace } X \mid b. b \in Ua\} \implies \text{openin } X c$
using *assms u(3)* **by** *blast*
show *?thesis*
by (*metis (mono-tags, lifting) 0 1 2 3 gdelta-in-Inter open-imp-gdelta-in*)
qed

lemma *g-delta-of-inj-open-map*:
assumes *open-map X Y f inj-on f (topspace X) gdelta-in X a*
shows *gdelta-in Y (f -' a)*
proof –
obtain Ua **where** u :

$Ua \neq \{\}$ countable $Ua \wedge b. b \in Ua \implies \text{openin } X \ b \ a = \bigcap Ua$
using *gdelta-inD*[*OF assms*(3)] **by** *metis*
then obtain j **where** $j \in Ua$ **by** *auto*
have $f' a = f' \bigcap Ua$ **by**(*simp add: u(4)*)
also have $\dots = \bigcap ((\cdot) f' Ua)$
using u *openin-subset* **by**(*auto intro!: image-INT*[*OF assms*(2) - $\langle j \in Ua \rangle$, *of id, simplified*])
also have $\dots = \bigcap \{f' u \mid u. u \in Ua\}$ **by** *auto*
finally have $0: f' a = \bigcap \{f' u \mid u. u \in Ua\}$.
have $1: \{f' u \mid u. u \in Ua\} \neq \{\}$
using $u(1)$ **by** *auto*
have $2: \text{countable } \{f' u \mid u. u \in Ua\}$
using $u(2)$ **by** (*simp add: Setcompr-eq-image*)
have $3: \bigwedge c. c \in \{f' u \mid u. u \in Ua\} \implies \text{openin } Y \ c$
using *assms*(1) $u(3)$ **by**(*auto simp: open-map-def*)
show *?thesis*
by (*metis* (*no-types*, *lifting*) 0 1 2 3 *gdelta-in-Inter open-imp-gdelta-in*)
qed

lemma *gdelta-in-prod*:

assumes *gdelta-in* $X \ A$ *gdelta-in* $Y \ B$
shows *gdelta-in* (*prod-topology* $X \ Y$) ($A \times B$)
proof –
obtain $Ua \ Ub$ **where** hu :
 $Ua \neq \{\}$ countable $Ua \wedge b. b \in Ua \implies \text{openin } X \ b \ A = \bigcap Ua$
 $Ub \neq \{\}$ countable $Ub \wedge b. b \in Ub \implies \text{openin } Y \ b \ B = \bigcap Ub$
by (*meson gdelta-inD assms*)
then have $0: A \times B = \bigcap \{a \times b \mid a \ b. a \in Ua \wedge b \in Ub\}$ **by** *blast*
have $1: \{a \times b \mid a \ b. a \in Ua \wedge b \in Ub\} \neq \{\}$
using $hu(1,5)$ **by** *auto*
have $2: \text{countable } \{a \times b \mid a \ b. a \in Ua \wedge b \in Ub\}$
proof –
have *countable* ($(\lambda(x, y). x \times y) ' (Ua \times Ub)$)
using $hu(2,6)$ **by**(*auto intro!: countable-image*[*of* $Ua \times Ub \ \lambda(x,y). x \times y$])
moreover have $\dots = \{a \times b \mid a \ b. a \in Ua \wedge b \in Ub\}$ **by** *auto*
ultimately show *?thesis* **by** *simp*
qed
have $3: \bigwedge c. c \in \{a \times b \mid a \ b. a \in Ua \wedge b \in Ub\} \implies \text{openin } (\text{prod-topology } X \ Y) \ c$
using $hu(3,7)$ **by**(*auto simp: openin-prod-Times-iff*)
show *?thesis*
by (*metis* (*no-types*, *lifting*) *gdelta-in-Inter open-imp-gdelta-in* 0 1 2 3)
qed

corollary *gdelta-in-prod1*:

assumes *gdelta-in* $X \ A$
shows *gdelta-in* (*prod-topology* $X \ Y$) ($A \times \text{topspace } Y$)
by(*auto intro!: gdelta-in-prod assms*)

corollary *gdelta-in-prod2*:
assumes *gdelta-in* $Y B$
shows *gdelta-in* (*prod-topology* $X Y$) (*topspace* $X \times B$)
by(*auto intro!*: *gdelta-in-prod assms*)

lemma *continuous-map-imp-closed-graph'*:
assumes *continuous-map* $X Y f$ *Hausdorff-space* Y
shows *closedin* (*prod-topology* $Y X$) ($(\lambda x. (f x, x))$ ' *topspace* X)
using *assms* *closed-map-def* *closed-map-paired-continuous-map-left* **by** *blast*

1.1.7 Continuous Maps on First Countable Topology

Generalized version of *Metric-space* $?M ?d \implies$ *eventually* $?P$ (*atin* (*Metric-space.mtopology* $?M ?d$) $?a$) = $(\forall \sigma. \text{range } \sigma \subseteq ?M - \{?a\} \wedge \text{limitin} (\text{Metric-space.mtopology } ?M ?d) \sigma ?a \text{ sequentially}) \longrightarrow (\forall_F n \text{ in sequentially. } ?P (\sigma n))$

lemma *eventually-atin-sequentially*:
assumes *first-countable* X
shows *eventually* P (*atin* $X a$) $\longleftrightarrow (\forall \sigma. \text{range } \sigma \subseteq \text{topspace } X - \{a\} \wedge \text{limitin } X \sigma a \text{ sequentially}) \longrightarrow \text{eventually } (\lambda n. P (\sigma n)) \text{ sequentially}$

proof *safe*

fix *an*

assume *h*:*eventually* P (*atin* $X a$) *range* *an* \subseteq *topspace* $X - \{a\}$ *limitin* $X \text{ an } a$ *sequentially*

then obtain U **where** $U: \text{openin } X U a \in U \forall x \in U - \{a\}. P x$

by (*auto simp: eventually-atin limitin-topspace*)

with $h(3)$ **obtain** N **where** $\forall n \geq N. \text{an } n \in U$

by (*meson limitin-sequentially*)

with $U(3)$ $h(2)$ **show** $\forall_F n \text{ in sequentially. } P (\text{an } n)$

unfolding *eventually-sequentially* **by** *blast*

next

assume *h*: $\forall \text{an. range } \text{an} \subseteq \text{topspace } X - \{a\} \wedge \text{limitin } X \text{ an } a \text{ sequentially} \longrightarrow (\forall_F n \text{ in sequentially. } P (\text{an } n))$

consider $a \notin \text{topspace } X \mid a \in \text{topspace } X$

by *blast*

then show *eventually* P (*atin* $X a$)

proof *cases*

assume $a: a \in \text{topspace } X$

from *a assms* **obtain** B' **where** $B': \text{countable } B' \wedge \forall V. V \in B' \implies \text{openin } X V \wedge U. \text{openin } X U \implies a \in U \implies (\exists V \in B'. a \in V \wedge V \subseteq U)$

by(*fastforce simp: first-countable-def*)

define B **where** $B \equiv \{V \in B'. a \in V\}$

have $B: \wedge V. V \in B \implies \text{openin } X V \text{ countable } B B \neq \{\}$ $\wedge U. \text{openin } X U \implies a \in U \implies (\exists V \in B. a \in V \wedge V \subseteq U)$

using $B' B'(3)[OF - a]$ **by**(*fastforce simp: B-def*)**+**

define An **where** $An \equiv (\lambda n. \bigcap_{i \leq n}. \text{from-nat-into } B i)$

have *a-in-An*: $a \in An n$ **for** n

by (*metis (no-types, lifting) An-def B-def B(3) INT-I from-nat-into mem-Collect-eq*)

have *openAn*: $\wedge n. \text{openin } X (An n)$

using B **by**(*auto simp: An-def from-nat-into[OF B(3)] openin-Inter*)


```

have deqseq-An:decseq An
  by(fastforce simp: decseq-def An-def)
have  $\exists U. \text{openin } X \ U \wedge a \in U \wedge \text{Ball } (U - \{a\}) \ P$ 
proof(rule ccontr)
  assume  $\nexists U. \text{openin } X \ U \wedge a \in U \wedge \text{Ball } (U - \{a\}) \ P$ 
  then have  $\bigwedge U. \text{openin } X \ U \implies a \in U \implies \exists x \in U - \{a\}. \neg P \ x$ 
    by blast
  hence  $\exists b \in An \ n - \{a\}. \neg P \ b$  for n
    using openAn a-in-An by auto
  then obtain an where an:  $\bigwedge n. an \ n \in An \ n \wedge n \neq a \wedge n. \neg P \ (an \ n)$ 
    by (metis Diff-iff singletonI)
  have limitin X an a sequentially
    unfolding limitin-sequentially
  proof safe
    fix U
    assume openin X U a ∈ U
    then obtain V where V:  $a \in V \ V \subseteq U \ V \in B$ 
      using B by meson
    then obtain N where V = from-nat-into B N
      by (metis B(2) from-nat-into-surj)
    hence  $\bigwedge n. n \geq N \implies an \ n \in V$ 
      using an(1) An-def by blast
    thus  $\exists N. \forall n \geq N. an \ n \in U$ 
      using V by blast
    qed fact
    hence  $1: \forall_F \ n \ \text{in} \ \text{sequentially}. P \ (an \ n)$ 
      using an(2) h an(1) openin-subset[OF openAn] by blast
    thus False
      using an(3) by simp
    qed
  thus ?thesis
    by(simp add: eventually-atin)
  qed(auto simp: eventually-atin)
qed

```

```

lemma continuous-map-iff-limit-seq:
  assumes first-countable X
  shows continuous-map X Y f  $\longleftrightarrow (\forall xn \ x. \text{limitin } X \ xn \ x \ \text{sequentially} \longrightarrow \text{limitin}$ 
Y  $(\lambda n. f \ (xn \ n)) \ (f \ x) \ \text{sequentially})$ 
  unfolding continuous-map-atin
proof safe
  fix xn x
  assume h:  $\forall x \in \text{topspace } X. \text{limitin } Y \ f \ (f \ x) \ (\text{atin } X \ x) \ \text{limitin } X \ xn \ x \ \text{sequentially}$ 
  then have limfx:  $\text{limitin } Y \ f \ (f \ x) \ (\text{atin } X \ x)$ 
    by(simp add: limitin-topospace)
  show  $\text{limitin } Y \ (\lambda n. f \ (xn \ n)) \ (f \ x) \ \text{sequentially}$ 
    unfolding limitin-sequentially
  proof safe
    fix U

```

```

assume  $U: \text{openin } Y \ U \ f \ x \in U$ 
then have  $h': \bigwedge \sigma. \text{range } \sigma \subseteq \text{topspace } X - \{x\} \implies x \in \text{topspace } X \implies \text{limitin } X \ \sigma \ x \text{ sequentially} \implies (\exists N. \forall n \geq N. f(\sigma \ n) \in U)$ 
using  $\text{limfx by}(\text{auto simp: limitin-def eventually-atin-sequentially}[OF \ \text{assms}(1)])$ 
eventually-sequentially)
show  $\exists N. \forall n \geq N. f(x \ n \ n) \in U$ 
proof(cases finite  $\{n. \ x \ n \ n \neq x\}$ )
  assume finite  $\{n. \ x \ n \ n \neq x\}$ 
  then obtain  $N$  where  $\bigwedge n. \ n \geq N \implies x \ n \ n = x$ 
    using infinite-nat-iff-unbounded-le by blast
  then show ?thesis
    using  $U$  by auto
next
assume inf:infinite  $\{n. \ x \ n \ n \neq x\}$ 
obtain  $n0$  where  $n0: \bigwedge n. \ n \geq n0 \implies x \ n \ n \in \text{topspace } X$ 
  by (meson  $h(2)$  limitin-sequentially openin-topspace)
have inf':infinite  $(\{n. \ x \ n \ n \neq x\} \cap \{n0..\})$ 
proof
  have  $1: (\{n. \ x \ n \ n \neq x\} \cap \{n0..\}) \cup (\{n. \ x \ n \ n \neq x\} \cap \{..<n0\}) = \{n. \ x \ n \ n \neq x\}$ 
    by auto
  assume finite  $(\{n. \ x \ n \ n \neq x\} \cap \{n0..\})$ 
  then have finite  $((\{n. \ x \ n \ n \neq x\} \cap \{n0..\}) \cup (\{n. \ x \ n \ n \neq x\} \cap \{..<n0\}))$ 
    by auto
  with inf show False
    unfolding  $1$  by blast
qed
define  $a$  where  $a \equiv \text{enumerate } (\{n. \ x \ n \ n \neq x\} \cap \{n0..\})$ 
have  $a: \text{strict-mono } a \ \text{range } a = (\{n. \ x \ n \ n \neq x\} \cap \{n0..\})$ 
  using range-enumerate[OF inf'] strict-mono-enumerate[OF inf']
  by(auto simp: a-def)
have  $\exists N. \forall n \geq N. f(x \ n \ (a \ n)) \in U$ 
  using limitin-subsequence[OF a(1) h(2)]  $a(2) \ n0$ 
  by(auto intro!: h' limitin-topspace[OF h(2)] simp: comp-def)
then obtain  $N$  where  $N: \bigwedge n. \ n \geq N \implies f(x \ n \ (a \ n)) \in U$ 
  by blast
show  $\exists N. \forall n \geq N. f(x \ n \ n) \in U$ 
proof(auto intro!: exI[where  $x=a \ N$ ])
  fix  $n$ 
  assume  $n: n \geq a \ N$ 
  show  $f(x \ n \ n) \in U$ 
  proof (cases  $x \ n \ n = x$ )
    assume  $x \ n \ n \neq x$ 
    moreover have  $n0 \leq n$ 
      using seq-suble[OF a(1), of N]  $n \ a(2)$ 
      by (metis Int-Collect atLeast-def dual-order.trans rangeI)
    ultimately obtain  $n1$  where  $n1: n = a \ n1$ 
    by (metis (mono-tags, lifting) Int-Collect atLeast-def imageE mem-Collect-eq
 $a(2)$ )

```

```

    have n1 ≥ N
      using strict-mono-less-eq[OF a(1), of N n1] n by (simp add: n1)
    thus ?thesis
      by (auto intro!: N simp: n1)
    qed (auto simp: U)
  qed
  qed
  qed (auto intro!: limitin-topospace limfx)
next
  fix x
  assume h: ∀ xn x. limitin X xn x sequentially → limitin Y (λn. f (xn n)) (f x)
  sequentially x ∈ topspace X
  then have f x ∈ topspace Y
    by (meson Abstract-Limits.limitin-const-iff limitin-topospace)
  thus limitin Y f (f x) (atin X x)
    using h by (auto simp: eventually-atin-sequentially[OF assms(1)] limitin-def )
  qed

```

1.1.8 Upper-Semicontinuous Functions

definition *upper-semicontinuous-map* :: [*'a topology, 'a ⇒ 'b :: linorder-topology*]
 \Rightarrow *bool* **where**
upper-semicontinuous-map $X f \longleftrightarrow (\forall a. \text{openin } X \{x \in \text{topspace } X. f x < a\})$

lemma *continuous-upper-semicontinuous*:
assumes *continuous-map* X (*euclidean* :: (*'b :: linorder-topology*) *topology*) f
shows *upper-semicontinuous-map* $X f$
unfolding *upper-semicontinuous-map-def*
proof *safe*
 fix $a :: 'b$
 have *: *openin euclidean* $U \implies \text{openin } X \{x \in \text{topspace } X. f x \in U\}$ **for** U
 using *assms* **by** (*simp add: continuous-map*)
 have *openin euclidean* $\{.. < a\}$ **by** *auto*
 with *[*of* $\{.. < a\}$] **show** *openin* $X \{x \in \text{topspace } X. f x < a\}$ **by** *auto*
 qed

lemma *upper-semicontinuous-map-iff-closed*:
upper-semicontinuous-map $X f \longleftrightarrow (\forall a. \text{closedin } X \{x \in \text{topspace } X. f x \geq a\})$
proof –
 have $\{x \in \text{topspace } X. f x < a\} = \text{topspace } X - \{x \in \text{topspace } X. f x \geq a\}$ **for**
 a
by *auto*
thus ?thesis
by (*simp add: closedin-def upper-semicontinuous-map-def*)
 qed

lemma *upper-semicontinuous-map-real-iff*:
fixes $f :: 'a \Rightarrow \text{real}$
shows *upper-semicontinuous-map* $X f \longleftrightarrow \text{upper-semicontinuous-map } X (\lambda x.$

```

ereal (f x))
  unfolding upper-semicontinuous-map-def
proof safe
  fix a :: ereal
  assume h:∀ a::real. openin X {x ∈ topspace X. f x < a}
  consider a = - ∞ | a = ∞ | a ≠ - ∞ ∧ a ≠ ∞ by auto
  then show openin X {x ∈ topspace X. ereal (f x) < a}
  proof cases
    case 3
    then have ereal (f x) < a ⟷ f x < real-of-ereal a for x
      by (metis ereal-less-eq(3) linorder-not-less real-of-ereal.elims)
    thus ?thesis
      using h by simp
  qed simp-all
next
  fix a :: real
  assume h:∀ a::ereal. openin X {x ∈ topspace X. ereal (f x) < a}
  then have openin X {x ∈ topspace X. ereal (f x) < ereal a}
    by blast
  moreover have ereal (f x) < real-of-ereal a ⟷ f x < a for x
    by auto
  ultimately show openin X {x ∈ topspace X. f x < a} by auto
qed

```

1.1.9 Lower-Semicontinuous Functions

definition *lower-semicontinuous-map* :: [*'a topology, 'a ⇒ 'b :: linorder-topology*]
 \Rightarrow *bool* **where**
lower-semicontinuous-map $X f \longleftrightarrow (\forall a. \text{openin } X \{x \in \text{topspace } X. a < f x\})$

lemma *continuous-lower-semicontinuous*:
assumes *continuous-map* X (*euclidean* :: (*'b :: linorder-topology*) *topology*) f
shows *lower-semicontinuous-map* $X f$
unfolding *lower-semicontinuous-map-def*

```

proof safe
  fix a :: 'b
  have *:openin euclidean U ⟹ openin X {x ∈ topspace X. f x ∈ U} for U
    using assms by(simp add: continuous-map)
  have openin euclidean {a<..} by auto
  with *[of {a<..}] show openin X {x ∈ topspace X. a < f x} by auto
qed

```

lemma *lower-semicontinuous-map-iff-closed*:
lower-semicontinuous-map $X f \longleftrightarrow (\forall a. \text{closedin } X \{x \in \text{topspace } X. f x \leq a\})$

```

proof -
  have {x ∈ topspace X. a < f x} = topspace X - {x ∈ topspace X. f x ≤ a} for
  a
  by auto
  thus ?thesis

```

by (*simp add: closedin-def lower-semicontinuous-map-def*)
qed

lemma *lower-semicontinuous-map-real-iff*:

fixes $f :: 'a \Rightarrow \text{real}$

shows $\text{lower-semicontinuous-map } X f \longleftrightarrow \text{lower-semicontinuous-map } X (\lambda x. \text{ereal } (f x))$

unfolding *lower-semicontinuous-map-def*

proof *safe*

fix $a :: \text{ereal}$

assume $h: \forall a :: \text{real}. \text{openin } X \{x \in \text{topspace } X. a < f x\}$

consider $a = -\infty \mid a = \infty \mid a \neq -\infty \wedge a \neq \infty$ **by** *auto*

then show $\text{openin } X \{x \in \text{topspace } X. a < \text{ereal } (f x)\}$

proof *cases*

case 3

then have $a < \text{ereal } (f x) \longleftrightarrow \text{real-of-ereal } a < f x$ **for** x

by (*metis ereal-less-eq(3) linorder-not-less real-of-ereal.elims*)

thus *?thesis*

using h **by** *simp*

qed *simp-all*

next

fix $a :: \text{real}$

assume $h: \forall a :: \text{ereal}. \text{openin } X \{x \in \text{topspace } X. a < \text{ereal } (f x)\}$

then have $\text{openin } X \{x \in \text{topspace } X. \text{ereal } (f x) > \text{ereal } a\}$

by *blast*

moreover have $\text{ereal } (f x) > \text{real-of-ereal } a \longleftrightarrow a < f x$ **for** x

by *auto*

ultimately show $\text{openin } X \{x \in \text{topspace } X. f x > a\}$ **by** *auto*

qed

1.2 Lemmas for Measure Theory

1.2.1 Lemmas for Measurable Sets

lemma *measurable-preserve-sigma-sets*:

assumes $\text{sets } M = \text{sigma-sets } \Omega S S \subseteq \text{Pow } \Omega$

$\bigwedge a. a \in S \implies f ' a \in \text{sets } N \text{ inj-on } f (\text{space } M) f ' \text{space } M \in \text{sets } N$

and $b \in \text{sets } M$

shows $f ' b \in \text{sets } N$

proof *–*

have $b \in \text{sigma-sets } \Omega S$

using *assms(1,6)* **by** *simp*

thus *?thesis*

proof *induction*

case (*Basic a*)

then show *?case* **by**(*rule assms(3)*)

next

case *Empty*

then show *?case* **by** *simp*

next

```

    case (Compl a)
    moreover have  $\Omega = \text{space } M$ 
    by (metis assms(1) assms(2) sets.sets-into-space sets.top sigma-sets-into-sp
sigma-sets-top subset-antisym)
    ultimately show ?case
    by (metis Diff-subset assms(2) assms(4) assms(5) inj-on-image-set-diff
sets.Diff sigma-sets-into-sp)
  next
  case (Union a)
  then show ?case
  by (simp add: image-UN)
qed
qed

```

```

inductive-set sigma-sets-cinter :: 'a set  $\Rightarrow$  'a set set  $\Rightarrow$  'a set set
for sp :: 'a set and A :: 'a set set
where
  Basic-c[intro, simp]:  $a \in A \implies a \in \text{sigma-sets-cinter } sp \ A$ 
  | Top-c[simp]:  $sp \in \text{sigma-sets-cinter } sp \ A$ 
  | Inter-c:  $(\bigwedge i::\text{nat. } a \ i \in \text{sigma-sets-cinter } sp \ A) \implies (\bigcap i. a \ i) \in \text{sigma-sets-cinter } sp \ A$ 
  | Union-c:  $(\bigwedge i::\text{nat. } a \ i \in \text{sigma-sets-cinter } sp \ A) \implies (\bigcup i. a \ i) \in \text{sigma-sets-cinter } sp \ A$ 

```

```

inductive-set sigma-sets-cinter-dunion :: 'a set  $\Rightarrow$  'a set set  $\Rightarrow$  'a set set
for sp :: 'a set and A :: 'a set set
where
  Basic-cd[intro, simp]:  $a \in A \implies a \in \text{sigma-sets-cinter-dunion } sp \ A$ 
  | Top-cd[simp]:  $sp \in \text{sigma-sets-cinter-dunion } sp \ A$ 
  | Inter-cd:  $(\bigwedge i::\text{nat. } a \ i \in \text{sigma-sets-cinter-dunion } sp \ A) \implies (\bigcap i. a \ i) \in \text{sigma-sets-cinter-dunion } sp \ A$ 
  | Union-cd:  $(\bigwedge i::\text{nat. } a \ i \in \text{sigma-sets-cinter-dunion } sp \ A) \implies \text{disjoint-family } a \implies (\bigcup i. a \ i) \in \text{sigma-sets-cinter-dunion } sp \ A$ 

```

lemma sigma-sets-cinter-dunion-subset: $\text{sigma-sets-cinter-dunion } sp \ A \subseteq \text{sigma-sets-cinter } sp \ A$

```

proof safe
  fix x
  assume  $x \in \text{sigma-sets-cinter-dunion } sp \ A$ 
  then show  $x \in \text{sigma-sets-cinter } sp \ A$ 
  by induction (auto intro!: Union-c Inter-c)
qed

```

```

lemma sigma-sets-cinter-into-sp:
  assumes  $A \subseteq \text{Pow } sp$   $x \in \text{sigma-sets-cinter } sp \ A$ 
  shows  $x \subseteq sp$ 
  using assms(2) by induction (use assms(1) subsetD in blast)+

```

lemma sigma-sets-cinter-dunion-into-sp:

```

assumes  $A \subseteq \text{Pow } sp$   $x \in \text{sigma-sets-cinter-dunion } sp$   $A$ 
shows  $x \subseteq sp$ 
using assms(2) by induction (use assms(1) subsetD in blast)+

lemma sigma-sets-cinter-int:
assumes  $a \in \text{sigma-sets-cinter } sp$   $A$   $b \in \text{sigma-sets-cinter } sp$   $A$ 
shows  $a \cap b \in \text{sigma-sets-cinter } sp$   $A$ 
proof -
  have  $1: a \cap b = (\bigcap i::nat. \text{if } i = 0 \text{ then } a \text{ else } b)$  by auto
  show ?thesis
    unfolding 1 by(rule Inter-c,use assms in auto)
qed

lemma sigma-sets-cinter-dunion-int:
assumes  $a \in \text{sigma-sets-cinter-dunion } sp$   $A$   $b \in \text{sigma-sets-cinter-dunion } sp$   $A$ 
shows  $a \cap b \in \text{sigma-sets-cinter-dunion } sp$   $A$ 
proof -
  have  $1: a \cap b = (\bigcap i::nat. \text{if } i = 0 \text{ then } a \text{ else } b)$  by auto
  show ?thesis
    unfolding 1 by(rule Inter-cd,use assms in auto)
qed

lemma sigma-sets-cinter-un:
assumes  $a \in \text{sigma-sets-cinter } sp$   $A$   $b \in \text{sigma-sets-cinter } sp$   $A$ 
shows  $a \cup b \in \text{sigma-sets-cinter } sp$   $A$ 
proof -
  have  $1: a \cup b = (\bigcup i::nat. \text{if } i = 0 \text{ then } a \text{ else } b)$  by auto
  show ?thesis
    unfolding 1 by(rule Union-c,use assms in auto)
qed

lemma sigma-sets-eq-cinter-dunion:
assumes metrizable-space  $X$ 
shows  $\text{sigma-sets } (\text{topspace } X) \{U. \text{openin } X \ U\} = \text{sigma-sets-cinter-dunion}$ 
 $(\text{topspace } X) \{U. \text{openin } X \ U\}$ 
proof safe
  fix  $a$ 
  interpret  $sa: \text{sigma-algebra } \text{topspace } X \ \text{sigma-sets } (\text{topspace } X) \{U. \text{openin } X$ 
 $U\}$ 
    by(auto intro!: sigma-algebra-sigma-sets openin-subset)
  assume  $a \in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\}$ 
  then show  $a \in \text{sigma-sets } (\text{topspace } X) \{U. \text{openin } X \ U\}$ 
    by induction auto
next
  have  $c: \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\} \subseteq \{U \in \text{sigma-sets-cinter-dunion}$ 
 $(\text{topspace } X) \{U. \text{openin } X \ U\}. \text{topspace } X - U \in \text{sigma-sets-cinter-dunion } (\text{topspace}$ 
 $X) \{U. \text{openin } X \ U\}\}$ 
proof
  fix  $a$ 

```

```

assume  $a: a \in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\}$ 
then show  $a \in \{U \in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\}. \text{topspace } X - U \in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\}\}$ 
proof induction
  case  $a:(\text{Basic-cd } a)$ 
  then have  $\text{gdelta-in } X \ (\text{topspace } X - a)$ 
    by (auto intro!: closed-imp-gdelta-in assms)
  from  $\text{gdelta-inD}'[\text{OF this}]$  obtain  $U$  where  $U:$ 
     $\bigwedge n :: \text{nat. openin } X \ (U \ n) \ \text{topspace } X - a = \bigcap \ (\text{range } U)$  by auto
  show ?case
    using  $a \ U(1)$  by (auto simp: U(2) intro!: Inter-cd)
next
  case Top-cd
  then show ?case by auto
next
  case  $ca:(\text{Inter-cd } a)$ 
  define  $b$  where  $b \equiv (\lambda n. (\text{topspace } X - a \ n) \cap (\bigcap i. \text{if } i < n \text{ then } a \ i \ \text{else } \text{topspace } X))$ 
  have  $bd:\text{disjoint-family } b$ 
    using nat-neq-iff by (fastforce simp: disjoint-family-on-def b-def)
  have  $\text{bin}: b \ i \in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\}$  for  $i$ 
    unfolding b-def
    apply (rule sigma-sets-cinter-dunion-int)
    using  $ca(2)[\text{of } i]$ 
    apply auto[1]
  apply (rule Inter-cd) using  $ca$  by auto
  have  $\text{bun}:\text{topspace } X - (\bigcap \ (\text{range } a)) = (\bigcup i. b \ i)$  (is ?lhs = ?rhs)
proof -
  { fix  $x$ 
    have  $x \in ?lhs \longleftrightarrow x \in \text{topspace } X \wedge x \in (\bigcup i. \text{topspace } X - a \ i)$ 
      by auto
    also have  $\dots \longleftrightarrow x \in \text{topspace } X \wedge (\exists n. x \in \text{topspace } X - a \ n)$ 
      by auto
    also have  $\dots \longleftrightarrow x \in \text{topspace } X \wedge (\exists n. x \in \text{topspace } X - a \ n \wedge (\forall i < n. x \in a \ i))$ 
  }
proof safe
  fix  $n$ 
  assume  $1: x \notin a \ n \ x \in \text{topspace } X$ 
  define  $N$  where  $N \equiv \text{Min } \{m. m \leq n \wedge x \notin a \ m\}$ 
  have  $N: x \notin a \ N \ N \leq n$ 
    using linorder-class.Min-in[of  $\{m. m \leq n \wedge x \notin a \ m\}$ ] 1
    by (auto simp: N-def)
  have  $N': x \in a \ i \ \text{if } i < N$  for  $i$ 
proof (rule ccontr)
  assume  $x \notin a \ i$ 
  then have  $N \leq i$ 
    using linorder-class.Min-le[of  $\{m. m \leq n \wedge x \notin a \ m\}$ ]  $i$  that N(2)
    by (auto simp: N-def)
  with that show False by auto

```



```

      qed
      show  $\exists n. x \in \text{topspace } X - a \ n \wedge (\forall i < n. x \in a \ i)$ 
        using  $N \ N'$  by(auto intro!: exI[where x=N] 1)
      qed auto
      also have ...  $\longleftrightarrow x \in ?rhs$ 
        by(auto simp: b-def)
      finally have  $x \in ?lhs \longleftrightarrow x \in ?rhs . \}$ 
      thus ?thesis by auto
    qed
    have ...  $\in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\}$ 
      by(rule Union-cd) (use bin bd in auto)
    thus ?case
      using Inter-cd[of a,OF ca(1)] by(auto simp: bun)
  next
    case ca:(Union-cd a)
    have  $\text{topspace } X - (\bigcup (\text{range } a)) = (\bigcap i. (\text{topspace } X - a \ i))$ 
      by simp
    have ...  $\in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\}$ 
      by(rule Inter-cd) (use ca in auto)
    then show ?case
      using Union-cd[of a,OF ca(1,2)] by auto
  qed
qed
fix a
assume  $a \in \text{sigma-sets } (\text{topspace } X) \{U. \text{openin } X \ U\}$ 
then show  $a \in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\}$ 
proof induction
  case a:(Union a)
  define b where  $b \equiv (\lambda n. a \ n \cap (\bigcap i. \text{if } i < n \text{ then } \text{topspace } X - a \ i \text{ else } \text{topspace } X))$ 
  have bd:disjoint-family b
    by(auto simp: disjoint-family-on-def b-def) (metis Diff-iff UnCI image-eqI
linorder-neqE-nat mem-Collect-eq)
  have bin:b i  $\in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \{U. \text{openin } X \ U\}$  for i
  unfolding b-def
  apply(rule sigma-sets-cinter-dunion-int)
  using a(2)[of i]
  apply auto[1]
  apply(rule Inter-cd) using c a by auto
  have bun:( $\bigcup i. a \ i = \bigcup i. b \ i$ ) (is ?lhs = ?rhs)
proof -
  {
    fix x
    have  $x \in ?lhs \longleftrightarrow x \in \text{topspace } X \wedge x \in ?lhs$ 
      using sigma-sets-cinter-dunion-into-sp[OF - a(2)]
      by (metis UN-iff subsetD subset-Pow-Union topspace-def)
    also have ...  $\longleftrightarrow x \in \text{topspace } X \wedge (\exists n. x \in a \ n)$  by auto
    also have ...  $\longleftrightarrow x \in \text{topspace } X \wedge (\exists n. x \in a \ n \wedge (\forall i < n. x \in \text{topspace } X - a \ i))$ 

```

```

proof safe
  fix n
  assume 1:  $x \in \text{topspace } X \ x \in a \ n$ 
  define N where  $N \equiv \text{Min } \{m. m \leq n \wedge x \in a \ m\}$ 
  have N:  $x \in a \ N \ N \leq n$ 
    using linorder-class.Min-in[of  $\{m. m \leq n \wedge x \in a \ m\}$ ] 1
    by(auto simp: N-def)
  have N':  $x \notin a \ i$  if  $i < N$  for i
  proof(rule ccontr)
    assume  $\neg x \notin a \ i$ 
    then have  $N \leq i$ 
      using linorder-class.Min-le[of  $\{m. m \leq n \wedge x \in a \ m\}$  i] that N(2)
      by(auto simp: N-def)
    with that show False by auto
  qed
  show  $\exists n. x \in a \ n \wedge (\forall i < n. x \in \text{topspace } X - a \ i)$ 
    using N N' 1 by(auto intro!: exI[where x=N])
qed auto
also have ...  $\longleftrightarrow x \in ?rhs$ 
proof safe
  fix m
  assume  $x \in b \ m$ 
  then show  $x \in \text{topspace } X \ \exists n. x \in a \ n \wedge (\forall i < n. x \in \text{topspace } X - a \ i)$ 
    by(auto intro!: exI[where x=m] simp: b-def)
  qed(auto simp: b-def)
  finally have  $x \in ?lhs \longleftrightarrow x \in ?rhs . \}$ 
thus ?thesis by auto
qed
have ...  $\in \text{sigma-sets-cinter-dunion } (\text{topspace } X) \ \{U. \text{openin } X \ U\}$ 
  by(rule Union-cd) (use bin bd in auto)
thus ?case
  by(auto simp: bun)
qed(use c in auto)
qed

lemma sigma-sets-eq-cinter:
  assumes metrizable-space X
  shows  $\text{sigma-sets } (\text{topspace } X) \ \{U. \text{openin } X \ U\} = \text{sigma-sets-cinter } (\text{topspace } X) \ \{U. \text{openin } X \ U\}$ 
proof safe
  fix a
  interpret sa: sigma-algebra topspace X sigma-sets (topspace X) {U. openin X U}
  by(auto intro!: sigma-algebra-sigma-sets openin-subset)
  assume  $a \in \text{sigma-sets-cinter } (\text{topspace } X) \ \{U. \text{openin } X \ U\}$ 
  then show  $a \in \text{sigma-sets } (\text{topspace } X) \ \{U. \text{openin } X \ U\}$ 
    by induction auto
qed (use sigma-sets-cinter-dunion-subset sigma-sets-eq-cinter-dunion[OF assms] in auto)

```

1.2.2 Measurable Isomorphisms

definition *measurable-isomorphic-map*::['a measure, 'b measure, 'a \Rightarrow 'b] \Rightarrow bool
where

measurable-isomorphic-map $M N f \longleftrightarrow$ *bij-betw* f (*space* M) (*space* N) $\wedge f \in M \rightarrow_M N \wedge$ *the-inv-into* (*space* M) $f \in N \rightarrow_M M$

lemma *measurable-isomorphic-map-sets-cong*:

assumes *sets* $M =$ *sets* M' *sets* $N =$ *sets* N'

shows *measurable-isomorphic-map* $M N f \longleftrightarrow$ *measurable-isomorphic-map* $M' N' f$

by (*simp add: measurable-isomorphic-map-def sets-eq-imp-space-eq*[*OF assms*(1)] *sets-eq-imp-space-eq*[*OF assms*(2)] *measurable-cong-sets*[*OF assms*] *measurable-cong-sets*[*OF assms*(2,1)])

lemma *measurable-isomorphic-map-surj*:

assumes *measurable-isomorphic-map* $M N f$

shows $f ' \text{space } M = \text{space } N$

using *assms* **by** (*auto simp: measurable-isomorphic-map-def bij-betw-def*)

lemma *measurable-isomorphic-mapI*:

assumes *bij-betw* f (*space* M) (*space* N) $f \in M \rightarrow_M N$ *the-inv-into* (*space* M) $f \in N \rightarrow_M M$

shows *measurable-isomorphic-map* $M N f$

using *assms* **by** (*simp add: measurable-isomorphic-map-def*)

lemma *measurable-isomorphic-map-byWitness*:

assumes $f \in M \rightarrow_M N$ $g \in N \rightarrow_M M \wedge x. x \in \text{space } M \Longrightarrow g (f x) = x \wedge x. x \in \text{space } N \Longrightarrow f (g x) = x$

shows *measurable-isomorphic-map* $M N f$

proof –

have *:*bij-betw* f (*space* M) (*space* N)

using *assms* **by** (*auto intro!: bij-betw-byWitness*[**where** $f'=g$] *dest:measurable-space*)

show ?*thesis*

proof (*rule measurable-isomorphic-mapI*)

have *the-inv-into* (*space* M) $f x = g x$ **if** $x \in \text{space } N$ **for** x

by (*metis* * *assms*(2) *assms*(4) *bij-betw-imp-inj-on measurable-space that the-inv-into-f-f*)

thus *the-inv-into* (*space* M) $f \in N \rightarrow_M M$

using *measurable-cong assms*(2) **by** *blast*

qed (*simp-all add: * assms*(1))

qed

lemma *measurable-isomorphic-map-restrict-space*:

assumes $f \in M \rightarrow_M N \wedge A. A \in \text{sets } M \Longrightarrow f ' A \in \text{sets } N$ *inj-on* f (*space* M)

shows *measurable-isomorphic-map* M (*restrict-space* N ($f ' \text{space } M$)) f

proof (*rule measurable-isomorphic-mapI*)

show *bij-betw* f (*space* M) (*space* (*restrict-space* N ($f ' \text{space } M$)))

by (*simp add: assms*(2,3) *inj-on-imp-bij-betw*)

next

```

show  $f \in M \rightarrow_M \text{restrict-space } N (f \text{ ' space } M)$ 
  by (simp add: assms(1) measurable-restrict-space2)
next
show  $\text{the-inv-into (space } M) f \in \text{restrict-space } N (f \text{ ' space } M) \rightarrow_M M$ 
proof(rule measurableI)
  show  $x \in \text{space (restrict-space } N (f \text{ ' space } M)) \implies \text{the-inv-into (space } M) f x$ 
   $\in \text{space } M$  for  $x$ 
    by (simp add: assms(2,3) the-inv-into-into)
  next
  fix  $A$ 
  assume  $A \in \text{sets } M$ 
  have  $\text{the-inv-into (space } M) f \text{ - ' } A \cap \text{space (restrict-space } N (f \text{ ' space } M)) =$ 
   $f \text{ ' } A$ 
    by (simp add: (A ∈ sets M) assms(2,3) sets.sets-into-space the-inv-into-vimage)
    also note assms(2)[OF (A ∈ sets M)]
  finally show  $\text{the-inv-into (space } M) f \text{ - ' } A \cap \text{space (restrict-space } N (f \text{ ' space } M)) \in \text{sets (restrict-space } N (f \text{ ' space } M))$ 
    by (simp add: assms(2) sets-restrict-space-iff)
  qed
qed

lemma measurable-isomorphic-mapD':
  assumes measurable-isomorphic-map M N f
  shows  $\bigwedge A. A \in \text{sets } M \implies f \text{ ' } A \in \text{sets } N \wedge f \in M \rightarrow_M N$ 
   $\exists g. \text{bij-betw } g \text{ (space } N) \text{ (space } M) \wedge g \in N \rightarrow_M M \wedge (\forall x \in \text{space } M. g (f$ 
   $x) = x) \wedge (\forall x \in \text{space } N. f (g x) = x) \wedge (\forall A \in \text{sets } N. g \text{ ' } A \in \text{sets } M)$ 
  proof -
  have  $h: \text{bij-betw } f \text{ (space } M) \text{ (space } N) \wedge f \in M \rightarrow_M N \wedge \text{the-inv-into (space } M) f \in$ 
   $N \rightarrow_M M$ 
    using assms by(simp-all add: measurable-isomorphic-map-def)
  show  $f \text{ ' } A \in \text{sets } N$  if  $A \in \text{sets } M$  for  $A$ 
  proof -
  have  $f \text{ ' } A = \text{the-inv-into (space } M) f \text{ - ' } A \cap \text{space } N$ 
  using the-inv-into-vimage[OF bij-betw-imp-inj-on[OF h(1)] sets.sets-into-space[OF
  that]]
  by(simp add: bij-betw-imp-surj-on[OF h(1)])
  also have  $\dots \in \text{sets } N$ 
  using that h(3) by auto
  finally show ?thesis .
  qed
show  $f \in M \rightarrow_M N$ 
  using assms by(simp add: measurable-isomorphic-map-def)

show  $\exists g. \text{bij-betw } g \text{ (space } N) \text{ (space } M) \wedge g \in N \rightarrow_M M \wedge (\forall x \in \text{space } M. g$ 
   $(f x) = x) \wedge (\forall x \in \text{space } N. f (g x) = x) \wedge (\forall A \in \text{sets } N. g \text{ ' } A \in \text{sets } M)$ 
  proof(rule exI[where x=the-inv-into (space M) f])
  have  $*: \text{the-inv-into (space } M) f \text{ ' } A \in \text{sets } M$  if  $A \in \text{sets } N$  for  $A$ 
  proof -
  have  $\bigwedge x. x \in \text{space } M \implies \text{the-inv-into (space } N) (\text{the-inv-into (space } M) f)$ 

```

$x = f x$
by (*metis* *bij-betw-imp-inj-on* *bij-betw-the-inv-into* $h(1)$ $h(2)$ *measurable-space* *the-inv-into-f-f*)
from *vimage-inter-cong*[*of space* $M - f A$, *OF this*] *the-inv-into-vimage*[*OF* *bij-betw-imp-inj-on*[*OF* *bij-betw-the-inv-into*[*OF* $h(1)$]]] *sets.sets-into-space*[*OF that*]
bij-betw-imp-surj-on[*OF* *bij-betw-the-inv-into*[*OF* $h(1)$]]] *measurable-sets*[*OF* $h(2)$ *that*]
show *?thesis*
by *fastforce*
qed
show *bij-betw* (*the-inv-into* (*space* M) f) (*space* N) (*space* M) \wedge *the-inv-into* (*space* M) $f \in N \rightarrow_M M \wedge (\forall x \in \text{space } M. \text{the-inv-into } (\text{space } M) f (f x) = x) \wedge (\forall x \in \text{space } N. f (\text{the-inv-into } (\text{space } M) f x) = x) \wedge (\forall A \in \text{sets } N. \text{the-inv-into } (\text{space } M) f ' A \in \text{sets } M)$
using *bij-betw-the-inv-into*[*OF* $h(1)$]
by (*meson* * *bij-betw-imp-inj-on* *f-the-inv-into-f-bij-betw* $h(1)$ $h(3)$ *the-inv-into-f-f*)
qed
qed

lemma *measurable-isomorphic-map-inv*:
assumes *measurable-isomorphic-map* $M N f$
shows *measurable-isomorphic-map* $N M$ (*the-inv-into* (*space* M) f)
using *assms*[*simplified measurable-isomorphic-map-def*]
by(*auto intro!*: *measurable-isomorphic-map-byWitness*[**where** $g=f$] *bij-betw-the-inv-into* *f-the-inv-into-f-bij-betw*[*of* f] *bij-betw-imp-inj-on* *the-inv-into-f-f*)

lemma *measurable-isomorphic-map-comp*:
assumes *measurable-isomorphic-map* $M N f$ **and** *measurable-isomorphic-map* $N L g$
shows *measurable-isomorphic-map* $M L$ ($g \circ f$)
proof –
obtain $f' g'$ **where**
 $[measurable]: f' \in N \rightarrow_M M$ **and** $hf: \bigwedge x. x \in \text{space } M \implies f' (f x) = x \wedge x \in \text{space } N \implies f (f' x) = x$
and $[measurable]: g' \in L \rightarrow_M N$ **and** $hg: \bigwedge x. x \in \text{space } N \implies g' (g x) = x \wedge x \in \text{space } L \implies g (g' x) = x$
using *measurable-isomorphic-mapD'*[*OF* *assms*(1)] *measurable-isomorphic-mapD'*[*OF* *assms*(2)] **by** *metis*
have $[measurable]: f \in M \rightarrow_M N$ $g \in N \rightarrow_M L$
using *assms* **by**(*auto simp: measurable-isomorphic-map-def*)
from $hf hg$ *measurable-space*[*OF* $\langle f \in M \rightarrow_M N \rangle$] *measurable-space*[*OF* $\langle g' \in L \rightarrow_M N \rangle$] **show** *?thesis*
by(*auto intro!*: *measurable-isomorphic-map-byWitness*[**where** $g=f' \circ g'$])
qed

definition *measurable-isomorphic*::[*'a measure, 'b measure*] \Rightarrow *bool* (**infixr** *measurable'-isomorphic* 50) **where**
 $M \text{ measurable-isomorphic } N \iff (\exists f. \text{measurable-isomorphic-map } M N f)$

lemma *measurable-isomorphic-sets-cong*:

assumes *sets M = sets M' sets N = sets N'*

shows *M measurable-isomorphic N \longleftrightarrow M' measurable-isomorphic N'*

using *measurable-isomorphic-map-sets-cong[OF assms]*

by(*auto simp: measurable-isomorphic-def*)

lemma *measurable-isomorphicD*:

assumes *M measurable-isomorphic N*

shows $\exists f g. f \in M \rightarrow_M N \wedge g \in N \rightarrow_M M \wedge (\forall x \in \text{space } M. g (f x) = x) \wedge$
 $(\forall y \in \text{space } N. f (g y) = y) \wedge (\forall A \in \text{sets } M. f ` A \in \text{sets } N) \wedge (\forall A \in \text{sets } N. g ` A$
 $\in \text{sets } M)$

using *assms measurable-isomorphic-mapD'[of M N]*

by (*metis (mono-tags, lifting) measurable-isomorphic-def*)

lemma *measurable-isomorphic-cardinality-eq*:

assumes *M measurable-isomorphic N*

shows *space M \approx space N*

by (*meson assms eqpoll-def measurable-isomorphic-def measurable-isomorphic-map-def*)

lemma *measurable-isomorphic-count-spaces*: *count-space A measurable-isomorphic*
count-space B \longleftrightarrow A \approx B

proof

assume *A \approx B*

then obtain *f where f: bij-betw f A B*

by(*auto simp: eqpoll-def*)

then show *count-space A measurable-isomorphic count-space B*

by(*auto simp: measurable-isomorphic-def measurable-isomorphic-map-def bij-betw-def*
the-inv-into-into intro!: exI[where x=f])

qed(*use measurable-isomorphic-cardinality-eq in fastforce*)

lemma *measurable-isomorphic-byWitness*:

assumes *f $\in M \rightarrow_M N \wedge x. x \in \text{space } M \implies g (f x) = x$*

and *g $\in N \rightarrow_M M \wedge y. y \in \text{space } N \implies f (g y) = y$*

shows *M measurable-isomorphic N*

by(*auto simp: measurable-isomorphic-def assms intro!: exI[where x = f] mea-*
surable-isomorphic-map-byWitness[where g=g])

lemma *measurable-isomorphic-refl*:

M measurable-isomorphic M

by(*auto intro!: measurable-isomorphic-byWitness[where f=id and g=id]*)

lemma *measurable-isomorphic-sym*:

assumes *M measurable-isomorphic N*

shows *N measurable-isomorphic M*

using *assms measurable-isomorphic-map-inv[of M N]*

by(*auto simp: measurable-isomorphic-def*)

lemma *measurable-isomorphic-trans*:

assumes *M measurable-isomorphic N and N measurable-isomorphic L*

shows M measurable-isomorphic L
using *assms* measurable-isomorphic-map-comp[of M N - L]
by(*auto simp: measurable-isomorphic-def*)

lemma *measurable-isomorphic-empty*:
assumes $space\ M = \{\}$ $space\ N = \{\}$
shows M measurable-isomorphic N
using *assms* **by**(*auto intro!: measurable-isomorphic-byWitness*[**where** $f=undefined$
and $g=undefined$] *simp: measurable-empty-iff*)

lemma *measurable-isomorphic-empty1*:
assumes $space\ M = \{\}$ M measurable-isomorphic N
shows $space\ N = \{\}$
using *measurable-isomorphicD*[*OF* *assms*(2)] **by**(*auto simp: measurable-empty-iff*[*OF*
assms(1)])

lemma *measurable-isomorphic-empty2*:
assumes $space\ N = \{\}$ M measurable-isomorphic N
shows $space\ M = \{\}$
using *measurable-isomorphic-sym*[*OF* *assms*(2)] *assms*(1)
by(*simp add: measurable-isomorphic-empty1*)

lemma *measurable-lift-product*:
assumes $\bigwedge i. i \in I \implies f\ i \in (M\ i) \rightarrow_M (N\ i)$
shows $(\lambda x\ i. \text{if } i \in I \text{ then } f\ i\ (x\ i) \text{ else } undefined) \in (\prod_M\ i \in I. M\ i) \rightarrow_M (\prod_{i \in I. N\ i})$
using *measurable-space*[*OF* *assms*]
by(*auto intro!: measurable-PiM-single' simp: assms measurable-PiM-component-rev*
space-PiM PiE-iff)

lemma *measurable-isomorphic-map-lift-product*:
assumes $\bigwedge i. i \in I \implies \text{measurable-isomorphic-map } (M\ i)\ (N\ i)\ (h\ i)$
shows *measurable-isomorphic-map* $(\prod_M\ i \in I. M\ i)\ (\prod_M\ i \in I. N\ i)\ (\lambda x\ i. \text{if } i \in I$
*then } h\ i\ (x\ i) \text{ else } undefined)
proof –
obtain h' **where**
 $\bigwedge i. i \in I \implies h'\ i \in (N\ i) \rightarrow_M (M\ i) \wedge i\ x. i \in I \implies x \in \text{space } (M\ i) \implies h'\ i$
 $(h\ i\ x) = x \wedge i\ x. i \in I \implies x \in \text{space } (N\ i) \implies h\ i\ (h'\ i\ x) = x$
using *measurable-isomorphic-mapD'*(3)[*OF* *assms*] **by** *metis*
thus *?thesis*
by(*auto intro!: measurable-isomorphic-map-byWitness*[*OF* *measurable-lift-product*[*of*
 $I\ h\ M\ N, OF\ \text{measurable-isomorphic-mapD}'(2)$][*OF* *assms*]] *measurable-lift-product*[*of*
 $I\ h'\ N\ M, OF\ \langle \bigwedge i. i \in I \implies h'\ i \in (N\ i) \rightarrow_M (M\ i) \rangle$]]
simp: space-PiM PiE-iff extensional-def)*

qed

lemma *measurable-isomorphic-lift-product*:
assumes $\bigwedge i. i \in I \implies (M\ i)$ measurable-isomorphic $(N\ i)$
shows $(\prod_M\ i \in I. M\ i)$ measurable-isomorphic $(\prod_M\ i \in I. N\ i)$

proof –

obtain h **where** $\bigwedge i. i \in I \implies \text{measurable-isomorphic-map } (M\ i) (N\ i) (h\ i)$
using assms **by** $(\text{auto simp: measurable-isomorphic-def})$ metis
thus $?thesis$
by $(\text{auto intro!: measurable-isomorphic-map-lift-product exI}[\text{where } x = \lambda x\ i. \text{ if } i \in I \text{ then } h\ i\ (x\ i) \text{ else undefined}] \text{simp: measurable-isomorphic-def})$
qed

<https://math24.net/cantor-schroder-bernstein-theorem.html>

lemma *Schroeder-Bernstein-measurable'*:

assumes $f' \text{ (space } M) \in \text{sets } N\ g' \text{ (space } N) \in \text{sets } M$
and $\text{measurable-isomorphic-map } M \text{ (restrict-space } N \text{ (} f' \text{ (space } M)))$ f **and**
 $\text{measurable-isomorphic-map } N \text{ (restrict-space } M \text{ (} g' \text{ (space } N)))$ g
shows $\exists h. \text{measurable-isomorphic-map } M\ N\ h$

proof –

have $\text{hset}:\bigwedge A. A \in \text{sets } M \implies f' \text{ (space } M) \in \text{sets } N$
 $\bigwedge A. A \in \text{sets } N \implies g' \text{ (space } N) \in \text{sets } M$
and $\text{hfg}[\text{measurable}]: f \in M \rightarrow_M \text{restrict-space } N \text{ (} f' \text{ (space } M))$
 $g \in N \rightarrow_M \text{restrict-space } M \text{ (} g' \text{ (space } N))$
using $\text{measurable-isomorphic-mapD}'(1,2)[\text{OF } \text{assms}(3)]$ $\text{measurable-isomorphic-mapD}'(1,2)[\text{OF } \text{assms}(4)]$ $\text{assms}(1,2)$
by auto
have $\text{hset2}:\bigwedge A. A \in \text{sets } M \implies f' \text{ (space } M) \in \text{sets } N$ $\bigwedge A. A \in \text{sets } N \implies g' \text{ (space } N) \in \text{sets } M$
and $\text{hfg2}[\text{measurable}]: f \in M \rightarrow_M N\ g \in N \rightarrow_M M$
using $\text{sets.Int-space-eq2}[\text{OF } \text{assms}(1)]$ $\text{sets.Int-space-eq2}[\text{OF } \text{assms}(2)]$ $\text{sets-restrict-space-iff}[\text{of } f' \text{ (space } M\ N)]$ $\text{sets-restrict-space-iff}[\text{of } g' \text{ (space } N\ M)]$ hset
 $\text{measurable-restrict-space2-iff}[\text{of } f\ M\ N]$ $\text{measurable-restrict-space2-iff}[\text{of } g\ N\ M]$ $\text{hfg } \text{assms}(1,2)$
by auto
have $\text{bij1}:\text{bij-betw } f \text{ (space } M) \text{ (} f' \text{ (space } M))$ $\text{bij-betw } g \text{ (space } N) \text{ (} g' \text{ (space } N))$
using $\text{assms}(3,4)$ **by** $(\text{auto simp: measurable-isomorphic-map-def space-restrict-space sets.Int-space-eq2}[\text{OF } \text{assms}(1)] \text{sets.Int-space-eq2}[\text{OF } \text{assms}(2)])$
obtain $f' g'$ **where**
 $\text{hfg1}[\text{measurable}]: f' \in \text{restrict-space } N \text{ (} f' \text{ (space } M)) \rightarrow_M M$ $g' \in \text{restrict-space } M \text{ (} g' \text{ (space } N)) \rightarrow_M N$
and $\text{hfg}':\bigwedge x. x \in \text{space } M \implies f' (f' x) = x$ $\bigwedge x. x \in f' \text{ (space } M) \implies f (f' x) = x$
 $\bigwedge x. x \in \text{space } N \implies g' (g' x) = x$ $\bigwedge x. x \in g' \text{ (space } N) \implies g (g' x) = x$
 $\text{bij-betw } f' \text{ (} f' \text{ (space } M)) \text{ (space } M)$ $\text{bij-betw } g' \text{ (} g' \text{ (space } N)) \text{ (space } N)$
using $\text{measurable-isomorphic-mapD}'(3)[\text{OF } \text{assms}(3)]$ $\text{measurable-isomorphic-mapD}'(3)[\text{OF } \text{assms}(4)]$ $\text{sets.Int-space-eq2}[\text{OF } \text{assms}(1)]$ $\text{sets.Int-space-eq2}[\text{OF } \text{assms}(2)]$
by $(\text{metis space-restrict-space})$

have $\text{hgfA}:(g \circ f)' \text{ (space } M) \in \text{sets } M$ **if** $A \in \text{sets } M$ **for** A

using $\text{hset2}(2)[\text{OF } \text{hset2}(1)[\text{OF } \text{that}]]$ **by** $(\text{simp add: image-comp})$
define An **where** $An \equiv (\lambda n. ((g \circ f)' \text{ (space } M - g' \text{ (space } N))) \text{ (} n))$
define A **where** $A \equiv (\bigcup n \in \text{UNIV}. An\ n)$
have $An\ n \in \text{sets } M$ **for** n


```

proof(induction n)
  case 0
  thus ?case
    using hset2[OF sets.top] by(simp add: An-def)
next
  case ih:(Suc n)
  have An (Suc n) = (g ∘ f) ' (An n)
    by(auto simp add: An-def)
  thus ?case
    using hgfA[OF ih] by simp
qed
hence Asets:A ∈ sets M
  by(simp add: A-def)
have Acompl:space M - A ⊆ g ' space N
proof -
  have space M - A ⊆ space M - An 0
    by(auto simp: A-def)
  also have ... ⊆ g ' space N
    by(auto simp: An-def)
  finally show ?thesis .
qed
define h where h ≡ (λx. if x ∈ A ∪ (- space M) then f x else g' x)
define h' where h' ≡ (λx. if x ∈ f ' A then f' x else g x)
have xinA-iff:x ∈ A ⟷ h x ∈ f ' A if x ∈ space M for x
proof
  assume h x ∈ f ' A
  show x ∈ A
  proof(rule ccontr)
    assume x ∉ A
    then have  $\bigwedge n. x \notin An\ n$ 
      by(auto simp: A-def)
    from this[of 0] have x ∈ g ' (space N)
      using that by(auto simp: An-def)
    have g' x ∈ f ' A
      using  $\langle h\ x \in f\ ' A \rangle$   $\langle x \notin A \rangle$ 
      by (simp add: h-def that)
    hence g (g' x) ∈ (g ∘ f) ' A
      by auto
    hence x ∈ (g ∘ f) ' A
      using  $\langle x \in g\ ' (space\ N) \rangle$  by (simp add: hgf'(4))
    then obtain n where x ∈ (g ∘ f) ' (An n)
      by(auto simp: A-def)
    hence x ∈ An (Suc n)
      by(auto simp: An-def)
    thus False
      using  $\langle \bigwedge n. x \notin An\ n \rangle$  by simp
  qed
qed(simp add: h-def)

```

show *?thesis*
proof(*intro exI*[**where** $x=h$] *measurable-isomorphic-map-byWitness*[**where** $g=h'$])
 have $\{x \in \text{space } M. x \in A \cup (- \text{space } M)\} \in \text{sets } M$
 using *sets.Int-space-eq2*[*OF Asets*] *Asets* **by** *simp*
 moreover **have** $f \in \text{restrict-space } M \{x. x \in A \cup - \text{space } M\} \rightarrow_M N$
 by (*simp add: measurable-restrict-space1*)
 moreover **have** $g' \in \text{restrict-space } M \{x. x \notin A \cup (- \text{space } M)\} \rightarrow_M N$
 proof –
 have *sets* (*restrict-space* (*restrict-space* M ($g' \text{ space } N$)) $\{x. x \notin A \cup - \text{space } M\}$) = *sets* (*restrict-space* M ($g' \text{ space } N \cap \{x. x \notin A \cup - \text{space } M\}$))
 by(*simp add: sets-restrict-restrict-space*)
 also **have** $\dots = \text{sets}$ (*restrict-space* M ($g' \text{ space } N \cap \{x. x \in \text{space } M - A\}$))
 by (*metis Compl-iff DiffE DiffI Un-iff*)
 also **have** $\dots = \text{sets}$ (*restrict-space* M $\{x. x \in \text{space } M - A\}$)
 by (*metis Acompl le-inf-iff mem-Collect-eq subsetI subset-antisym*)
 also **have** $\dots = \text{sets}$ (*restrict-space* M $\{x. x \notin A \cup (- \text{space } M)\}$)
 by (*metis Compl-iff DiffE DiffI Un-iff*)
 finally **have** *sets* (*restrict-space* (*restrict-space* M ($g' \text{ space } N$)) $\{x. x \notin A \cup - \text{space } M\}$) = *sets* (*restrict-space* M $\{x. x \notin A \cup - \text{space } M\}$) .
 from *measurable-cong-sets*[*OF this refl*] *measurable-restrict-space1*[*OF hfg1'(2), of*
 $\{x. x \notin A \cup - \text{space } M\}$]
 show *?thesis* **by** *auto*
 qed
 ultimately **show** $h \in M \rightarrow_M N$
 by(*simp add: h-def measurable-If-restrict-space-iff*)
next
 have $\{x \in \text{space } N. x \in f' A\} \in \text{sets } N$
 using *sets.Int-space-eq2*[*OF hset2(1)*][*OF Asets*] *hset2(1)*[*OF Asets*] **by** *simp*
 moreover **have** $f' \in \text{restrict-space } N \{x. x \in f' A\} \rightarrow_M M$
 proof –
 have *sets* (*restrict-space* (*restrict-space* N ($f' \text{ space } M$)) $\{x. x \in f' A\}$) =
sets (*restrict-space* N ($f' \text{ space } M \cap \{x. x \in f' A\}$))
 by(*simp add: sets-restrict-restrict-space*)
 also **have** $\dots = \text{sets}$ (*restrict-space* N $\{x. x \in f' A\}$)
 proof –
 have $f' \text{ space } M \cap \{x. x \in f' A\} = \{x. x \in f' A\}$
 using *sets.sets-into-space*[*OF Asets*] **by** *auto*
 thus *?thesis* **by** *simp*
 qed
 finally **have** *sets* (*restrict-space* (*restrict-space* N ($f' \text{ space } M$)) $\{x. x \in f' A\}$) = *sets* (*restrict-space* N $\{x. x \in f' A\}$) .
 from *measurable-cong-sets*[*OF this refl*] *measurable-restrict-space1*[*OF hfg1'(1), of*
 $\{x. x \in f' A\}$]
 show *?thesis* **by** *auto*
 qed
 moreover **have** $g \in \text{restrict-space } N \{x. x \notin f' A\} \rightarrow_M M$
 by (*simp add: measurable-restrict-space1*)
 ultimately **show** $h' \in N \rightarrow_M M$
 by(*simp add: h'-def measurable-If-restrict-space-iff*)

```

next
  fix x
  assume x ∈ space M
  then consider x ∈ A | x ∈ space M - A by auto
  thus h' (h x) = x
  proof cases
    case xa:2
      hence h x ∉ f ' A
      using ⟨x ∈ space M⟩ xinA-iff by blast
      thus ?thesis
      using Acompl hfg'(4) xa by(auto simp add: h-def h'-def)
  qed(simp add: h-def h'-def ⟨x ∈ space M⟩ hfg'(1))
next
  fix x
  assume x ∈ space N
  then consider x ∈ f ' A | x ∈ space N - f ' A by auto
  thus h (h' x) = x
  proof cases
    case hx:1
      hence x ∈ f ' (space M)
      using image-mono[OF sets.sets-into-space[OF Asets],of f] by auto
      have h' x = f' x
      using hx by(simp add: h'-def)
      also have ... ∈ A
      using hx sets.sets-into-space[OF Asets] hfg'(1) by auto
      finally show ?thesis
      using hfg'(2)[OF ⟨x ∈ f ' (space M)⟩] hx by(auto simp: h-def h'-def)
    case hx:2
      then have h' x = g x
      by(simp add: h'-def)
      also have ... ∉ A
      proof(rule ccontr)
        assume ¬ g x ∉ A
        then have g x ∈ A by simp
        then obtain n where hg:g x ∈ An n by(auto simp: A-def)
        hence 0 < n using hx by(auto simp: An-def)
        then obtain n' where [simp]:n = Suc n'
        using not0-implies-Suc by blast
        then have g x ∈ g ' f ' An n'
        using hg by(auto simp: An-def)
        hence x ∈ f ' An n'
        using inj-on-image-mem-iff[OF bij-betw-imp-inj-on[OF bij1(2)] ⟨x ∈ space
N⟩,of f ' An n']
        sets.sets-into-space[OF ⟨An n' ∈ sets M⟩] measurable-space[OF hfg2(1)]
      by auto
      also have ... ⊆ f ' A
      by(auto simp: A-def)
      finally show False

```

```

    using hx by simp
  qed
  finally show ?thesis
    using hx hfg'(3)[OF ⟨x ∈ space N⟩] measurable-space[OF hfg2(2) ⟨x ∈ space
N⟩]
    by(auto simp: h-def h'-def)
  qed
  qed
  qed

```

lemma *Schroeder-Bernstein-measurable:*

```

  assumes f ∈ M →M N ∧ A. A ∈ sets M ⇒ f ' A ∈ sets N inj-on f (space M)
    and g ∈ N →M M ∧ A. A ∈ sets N ⇒ g ' A ∈ sets M inj-on g (space N)
  shows ∃ h. measurable-isomorphic-map M N h
  using Schroeder-Bernstein-measurable'[OF assms(2)[OF sets.top] assms(5)[OF
sets.top] measurable-isomorphic-map-restrict-space[OF assms(1-3)] measurable-isomorphic-map-restrict-space
assms(4-6)]
  by simp

```

lemma *measurable-isomorphic-from-embeddings:*

```

  assumes M measurable-isomorphic (restrict-space N B) N measurable-isomorphic
(restrict-space M A)
    and A ∈ sets M B ∈ sets N
  shows M measurable-isomorphic N

```

proof –

```

  obtain f g where fg:measurable-isomorphic-map M (restrict-space N B) f mea-
surable-isomorphic-map N (restrict-space M A) g
    using assms(1,2) by(auto simp: measurable-isomorphic-def)
  have [simp]:f ' space M = B g ' space N = A
    using measurable-isomorphic-map-surj[OF fg(1)] measurable-isomorphic-map-surj[OF
fg(2)] sets.sets-into-space[OF assms(3)] sets.sets-into-space[OF assms(4)]
    by(auto simp: space-restrict-space)
  obtain h where measurable-isomorphic-map M N h
    using Schroeder-Bernstein-measurable'[of f M N g] assms(3,4) fg by auto
  thus ?thesis
    by(auto simp: measurable-isomorphic-def)
  qed

```

lemma *measurable-isomorphic-antisym:*

```

  assumes B measurable-isomorphic (restrict-space C c) A measurable-isomorphic
(restrict-space B b)
    and c ∈ sets C b ∈ sets B C measurable-isomorphic A
  shows C measurable-isomorphic B
  by(rule measurable-isomorphic-from-embeddings[OF measurable-isomorphic-trans[OF
assms(5,2)] assms(1) assms(3,4)])

```

lemma *countable-infinite-isomorphisc-to-nat-index:*

```

  assumes countable I and infinite I
  shows (ΠM x∈I. M) measurable-isomorphic (ΠM (x::nat)∈UNIV. M)

```

proof(*rule measurable-isomorphic-byWitness*[**where** $f = \lambda x n. x$ (*from-nat-into* I n) **and** $g = \lambda x. \lambda i \in I. x$ (*to-nat-on* I i)])
show $(\lambda x n. x$ (*from-nat-into* I n)) \in $(\prod_M x \in I. M) \rightarrow_M (\prod_M (x :: \text{nat}) \in \text{UNIV}. M)$
by(*auto intro!*: *measurable-PiM-single'* *measurable-component-singleton*[*OF from-nat-into*[*OF infinite-imp-nonempty*[*OF assms*(2)]]])
(*simp add*: *PiE-iff infinite-imp-nonempty space-PiM from-nat-into*[*OF infinite-imp-nonempty*[*OF assms*(2)]]])
next
show $(\lambda x. \lambda i \in I. x$ (*to-nat-on* I i)) \in $(\prod_M (x :: \text{nat}) \in \text{UNIV}. M) \rightarrow_M (\prod_M x \in I. M)$
by(*auto intro!*: *measurable-PiM-single'*)
next
show $x \in \text{space } (\prod_M x \in I. M) \implies (\lambda i \in I. x$ (*from-nat-into* I (*to-nat-on* I i))) = x **for** x
by (*simp add*: *assms*(1) *restrict-ext space-PiM*)
next
show $y \in \text{space } (P_{iM} \text{ UNIV } (\lambda x. M)) \implies (\lambda n. (\lambda i \in I. y$ (*to-nat-on* I i)) (*from-nat-into* I n)) = y **for** y
by (*simp add*: *assms*(1) *assms*(2) *from-nat-into infinite-imp-nonempty*)
qed

lemma *PiM-PiM-isomorphic-to-PiM*:

$(\prod_M i \in I. \prod_M j \in J. M i j)$ *measurable-isomorphic* $(\prod_M (i,j) \in I \times J. M i j)$
proof(*rule measurable-isomorphic-byWitness*[**where** $f = \lambda x (i,j). \text{if } (i,j) \in I \times J$ then $x i j$ else *undefined* **and** $g = \lambda x i j. \text{if } i \notin I$ then *undefined* j else *if* $j \notin J$ then *undefined* else $x (i,j)$])
have [*simp*]: $(\lambda \omega. \omega a b) \in (\prod_M i \in I. \prod_M j \in J. M i j) \rightarrow_M M a b$ **if** $a \in I b \in J$ **for** $a b$
using *measurable-component-singleton*[*OF that*(1),*of* $\lambda i. \prod_M j \in J. M i j$] *measurable-component-singleton*[*OF that*(2),*of* $M a$]
by *auto*
show $(\lambda x (i, j). \text{if } (i, j) \in I \times J$ then $x i j$ else *undefined*) \in $(\prod_M i \in I. \prod_M j \in J. M i j) \rightarrow_M (\prod_M (i,j) \in I \times J. M i j)$
apply(*rule measurable-PiM-single'*)
apply *auto*[1]
apply(*auto simp*: *PiE-def Pi-def space-PiM extensional-def; meson*)
done
next
have [*simp*]: $(\lambda \omega. \omega (i, j)) \in P_{iM} (I \times J) (\lambda (i, j). M i j) \rightarrow_M M i j$ **if** $i \in I j \in J$ **for** $i j$
using *measurable-component-singleton*[*of* $(i,j) I \times J \lambda (i, j). M i j$] **that** **by** *auto*
show $(\lambda x i j. \text{if } i \notin I$ then *undefined* j else *if* $j \notin J$ then *undefined* else $x (i, j)) \in (\prod_M (i,j) \in I \times J. M i j) \rightarrow_M (\prod_M i \in I. \prod_M j \in J. M i j)$
by(*auto intro!*: *measurable-PiM-single'*) (*simp-all add*: *PiE-iff space-PiM extensional-def*)
next
show $x \in \text{space } (\prod_M i \in I. \prod_M j \in J. M i j) \implies (\lambda i j. \text{if } i \notin I$ then *undefined* j

else if $j \notin J$ then undefined else case (i, j) of $(i, j) \Rightarrow$ if $(i, j) \in I \times J$ then $x \ i \ j$ else undefined) = x **for** x

by standard+ (auto simp: space-PiM PiE-def Pi-def extensional-def)

next

show $y \in \text{space } (\prod_{M (i,j) \in I \times J. M \ i \ j} \implies (\lambda(i, j). \text{if } (i, j) \in I \times J \text{ then if } i \notin I \text{ then undefined } j \text{ else if } j \notin J \text{ then undefined else } y \ (i, j) \text{ else undefined}) = y$ **for** y

by standard+ (auto simp: space-PiM PiE-def Pi-def extensional-def)

qed

lemma measurable-isomorphic-map-sigma-sets:

assumes sets $M = \text{sigma-sets } (\text{space } M) \ U$ measurable-isomorphic-map $M \ N \ f$

shows sets $N = \text{sigma-sets } (\text{space } N) \ ((\cdot) \ f \ U)$

proof –

from measurable-isomorphic-mapD'[OF assms(2)]

obtain g **where** $h: \bigwedge A. A \in \text{sets } M \implies f \ U \ A \in \text{sets } N \ f \in M \rightarrow_M N \text{ bij-betw } g$ ($\text{space } N$) ($\text{space } M$) $g \in N \rightarrow_M M \bigwedge x. x \in \text{space } M \implies g \ (f \ x) = x \bigwedge x. x \in \text{space } N \implies f \ (g \ x) = x \bigwedge A. A \in \text{sets } N \implies g \ U \ A \in \text{sets } M$

by metis

interpret $s: \text{sigma-algebra } \text{space } N \ \text{sigma-sets } (\text{space } N) \ ((\cdot) \ f \ U)$

by (auto intro!: sigma-algebra-sigma-sets) (metis assms(1) h(2) measurable-space sets.sets-into-space sigma-sets-superset-generator subsetD)

show ?thesis

proof safe

fix x

assume $x \in \text{sets } N$

from h(7)[OF this] assms(1)

have $g \ U \ x \in \text{sigma-sets } (\text{space } M) \ U$ **by** simp

hence $f \ U \ (g \ U \ x) \in \text{sigma-sets } (\text{space } N) \ ((\cdot) \ f \ U)$

proof induction

case $h:(\text{Compl } a)$

have $f \ U \ (\text{space } M - a) = f \ U \ (\text{space } M) - f \ U \ a$

by (rule inj-on-image-set-diff[**where** $C = \text{space } M$], insert assms h) (auto simp: measurable-isomorphic-map-def bij-betw-def sets.sets-into-space)

with h **show** ?case

by (metis assms(2) measurable-isomorphic-map-surj s.Diff s.top)

qed (auto simp: image-UN)

moreover **have** $f \ U \ (g \ U \ x) = x$

using sets.sets-into-space[OF $\langle x \in \text{sets } N \rangle$] h(6) **by** (fastforce simp: image-def)

ultimately **show** $x \in \text{sigma-sets } (\text{space } N) \ ((\cdot) \ f \ U)$ **by** simp

next

interpret $s': \text{sigma-algebra } \text{space } M \ \text{sigma-sets } (\text{space } M) \ U$

by (simp add: assms(1)[symmetric] sets.sigma-algebra-axioms)

have $1: \bigwedge x. x \in U \implies x \subseteq \text{space } M$

by (simp add: s'.sets-into-space)

fix x

assume $\text{assm}: x \in \text{sigma-sets } (\text{space } N) \ ((\cdot) \ f \ U)$

then **show** $x \in \text{sets } N$

by induction (auto simp: assms(1) h(1))

qed
qed

1.2.3 Borel Spaces Generated from Abstract Topologies

definition *borel-of* :: 'a topology \Rightarrow 'a measure **where**
borel-of $X \equiv \text{sigma} (\text{topspace } X) \{U. \text{openin } X U\}$

lemma *emeasure-borel-of*: *emeasure* (*borel-of* X) $A = 0$
by (*simp add: borel-of-def emeasure-sigma*)

lemma *borel-of-euclidean*: *borel-of euclidean* = *borel*
by(*simp add: borel-of-def borel-def*)

lemma *space-borel-of*: *space* (*borel-of* X) = *topspace* X
by(*simp add: space-measure-of-conv borel-of-def*)

lemma *sets-borel-of*: *sets* (*borel-of* X) = *sigma-sets* (*topspace* X) $\{U. \text{openin } X U\}$
by (*simp add: subset-Pow-Union topspace-def borel-of-def*)

lemma *sets-borel-of-closed*: *sets* (*borel-of* X) = *sigma-sets* (*topspace* X) $\{U. \text{closedin } X U\}$

unfolding *sets-borel-of*

proof(*safe intro!*: *sigma-sets-eqI*)

fix a

assume $a:\text{openin } X a$

have *topspace* $X - (\text{topspace } X - a) \in \text{sigma-sets} (\text{topspace } X) \{U. \text{closedin } X U\}$

by(*rule sigma-sets.Compl*) (*use a in auto*)

thus $a \in \text{sigma-sets} (\text{topspace } X) \{U. \text{closedin } X U\}$

using *openin-subset[OF a]* **by** (*simp add: Diff-Diff-Int inf.absorb-iff2*)

next

fix b

assume $b:\text{closedin } X b$

have *topspace* $X - (\text{topspace } X - b) \in \text{sigma-sets} (\text{topspace } X) \{U. \text{openin } X U\}$

by(*rule sigma-sets.Compl*) (*use b in auto*)

thus $b \in \text{sigma-sets} (\text{topspace } X) \{U. \text{openin } X U\}$

using *closedin-subset[OF b]* **by** (*simp add: Diff-Diff-Int inf.absorb-iff2*)

qed

lemma *borel-of-open*:

assumes *openin* $X U$

shows $U \in \text{sets} (\text{borel-of } X)$

using *assms* **by** (*simp add: subset-Pow-Union topspace-def borel-of-def*)

lemma *borel-of-closed*:

assumes *closedin* $X U$

shows $U \in \text{sets (borel-of } X)$
using *assms sigma-sets.Compl[of topspace X – U topspace X]*
by (*simp add: closedin-def double-diff sets-borel-of*)

lemma(in *Metric-space*) *nbh-sets[measurable]*: $(\bigcup a \in A. \text{mball } a \ e) \in \text{sets (borel-of } mtopology)$
by(*auto intro!: borel-of-open openin-clauses(3)*)

lemma *borel-of-gdelta-in*:
assumes *gdelta-in X U*
shows $U \in \text{sets (borel-of } X)$
using *gdelta-inD[OF assms] borel-of-open*
by(*auto intro!: sets.countable-INT[of - id,simplified]*)

lemma *borel-of-subtopology*:
borel-of (subtopology X U) = restrict-space (borel-of X) U
proof(*rule measure-eqI*)
show $\text{sets (borel-of (subtopology X U))} = \text{sets (restrict-space (borel-of X) U)}$
unfolding *restrict-space-eq-vimage-algebra' sets-vimage-algebra sets-borel-of topspace-subtopology space-borel-of Int-commute[of U]*
proof(*rule sigma-sets-eqI*)
fix *a*
assume $a \in \text{Collect (openin (subtopology X U))}$
then obtain *T* **where** *openin X T a = T ∩ U*
by(*auto simp: openin-subtopology*)
show $a \in \text{sigma-sets (topspace X ∩ U) } \{(\lambda x. x) - ' A \cap (\text{topspace X} \cap U) \mid A. A \in \text{sigma-sets (topspace X) (Collect (openin X))}\}$
using *openin-subset[OF ‹openin X T›] ‹a = T ∩ U›* **by**(*auto intro!: exI[where x=T] ‹openin X T›*)
next
fix *b*
assume $b \in \{(\lambda x. x) - ' A \cap (\text{topspace X} \cap U) \mid A. A \in \text{sigma-sets (topspace X) (Collect (openin X))}\}$
then obtain *T* **where** *ht:b = T ∩ (topspace X ∩ U) T ∈ sigma-sets (topspace X) (Collect (openin X))*
by *auto*
hence $b = T \cap U$
proof –
have $T \subseteq \text{topspace X}$
by(*rule sigma-sets-into-sp[OF - ht(2)] (simp add: subset-Pow-Union topspace-def)*)
thus *?thesis*
by(*auto simp: ht(1)*)
qed
with *ht(2)* **show** $b \in \text{sigma-sets (topspace X ∩ U) (Collect (openin (subtopology X U)))}$
proof(*induction arbitrary: b U*)
case (*Basic a*)
then show *?case*


```

    by(auto simp: openin-subtopology)
  next
    case Empty
    then show ?case by simp
  next
    case ih:(Compl a)
    then show ?case
      by (simp add: Diff-Int-distrib2 sigma-sets.Compl)
  next
    case (Union a)
    then show ?case
      by (metis UN-extend-simps(4) sigma-sets.Union)
  qed
qed
qed(simp add: emeasure-borel-of restrict-space-def emeasure-measure-of-conv)

lemma sets-borel-of-discrete-topology: sets (borel-of (discrete-topology I)) = sets
(count-space I)
  by (metis Pow-UNIV UNIV-eq-I borel-of-open borel-of-subtopology inf.absorb-iff2
openin-discrete-topology sets-count-space sets-restrict-space sets-restrict-space-count-space
subtopology-discrete-topology top-greatest)

lemma continuous-map-measurable:
  assumes continuous-map X Y f
  shows  $f \in \text{borel-of } X \rightarrow_M \text{borel-of } Y$ 
proof(rule measurable-sigma-sets[OF sets-borel-of[of Y]])
  show  $\{U. \text{openin } Y U\} \subseteq \text{Pow } (\text{topspace } Y)$ 
    by (simp add: subset-Pow-Union topspace-def)
  next
  show  $f \in \text{space } (\text{borel-of } X) \rightarrow \text{topspace } Y$ 
    using continuous-map-image-subset-topospace[OF assms]
    by(auto simp: space-borel-of)
  next
  fix U
  assume  $U \in \{U. \text{openin } Y U\}$ 
  then have  $\text{openin } X (f -' U \cap \text{topspace } X)$ 
    using continuous-map[of X Y f] assms by auto
  thus  $f -' U \cap \text{space } (\text{borel-of } X) \in \text{sets } (\text{borel-of } X)$ 
    by(simp add: space-borel-of sets-borel-of)
  qed

lemma upper-semicontinuous-map-measurable:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{linorder-topology, second-countable-topology}\}$ 
  assumes upper-semicontinuous-map X f
  shows  $f \in \text{borel-measurable } (\text{borel-of } X)$ 
  using assms by(auto intro!: borel-measurableI-less borel-of-open simp: space-borel-of
upper-semicontinuous-map-def)

lemma lower-semicontinuous-map-measurable:

```

fixes $f :: 'a \Rightarrow 'b :: \{\text{linorder-topology, second-countable-topology}\}$
assumes *lower-semicontinuous-map* $X f$
shows $f \in \text{borel-measurable (borel-of } X)$
using *assms* **by**(*auto intro!*: *borel-measurableI-greater borel-of-open simp: space-borel-of lower-semicontinuous-map-def*)

lemma *open-map-preserves-sets*:

assumes *open-map* $S T f \text{ inj-on } f \text{ (topspace } S) A \in \text{sets (borel-of } S)$
shows $f ' A \in \text{sets (borel-of } T)$
using *assms*(3)[*simplified sets-borel-of*]
proof(*induction*)
case (*Basic a*)
with *assms*(1) **show** ?*case*
by(*auto simp: sets-borel-of open-map-def*)
next
case *Empty*
show ?*case* **by** *simp*
next
case (*Compl a*)
moreover **have** $f ' (\text{topspace } S - a) = f ' (\text{topspace } S) - f ' a$
by (*metis Diff-subset assms*(2) *calculation*(1) *inj-on-image-set-diff sigma-sets-into-sp subset-Pow-Union topspace-def*)
moreover **have** $f ' (\text{topspace } S) \in \text{sets (borel-of } T)$
by (*meson assms*(1) *borel-of-open open-map-def openin-topspace*)
ultimately **show** ?*case*
by *auto*
next
case (*Union a*)
then **show** ?*case*
by (*simp add: image-UN*)
qed

lemma *open-map-preserves-sets'*:

assumes *open-map* $S \text{ (subtopology } T \text{ (} f ' (\text{topspace } S))) f \text{ inj-on } f \text{ (topspace } S)$
 $f ' (\text{topspace } S) \in \text{sets (borel-of } T) A \in \text{sets (borel-of } S)$
shows $f ' A \in \text{sets (borel-of } T)$
using *assms*(4)[*simplified sets-borel-of*]
proof(*induction*)
case (*Basic a*)
then **have** *openin* (*subtopology* $T \text{ (} f ' (\text{topspace } S))) \text{ (} f ' a$)
using *assms*(1) **by**(*auto simp: open-map-def*)
hence $f ' a \in \text{sets (borel-of (subtopology } T \text{ (} f ' (\text{topspace } S))))$
by(*simp add: sets-borel-of*)
hence $f ' a \in \text{sets (restrict-space (borel-of } T) \text{ (} f ' (\text{topspace } S)))$
by(*simp add: borel-of-subtopology*)
thus ?*case*
by (*metis sets-restrict-space-iff assms*(3) *sets.Int-space-eq2*)
next
case *Empty*

```

show ?case by simp
next
  case (Compl a)
  moreover have  $f'(\text{topspace } S - a) = f'(\text{topspace } S) - f' a$ 
  by (metis Diff-subset assms(2) calculation(1) inj-on-image-set-diff sigma-sets-into-sp
subset-Pow-Union topspace-def)
  ultimately show ?case
  using assms(3) by auto
next
  case (Union a)
  then show ?case
  by (simp add: image-UN)
qed

```

Abstract topology version of $\text{open} = \text{generate-topology } ?X \implies \text{borel} = \text{sigma UNIV } ?X$.

```

lemma borel-of-second-countable':
  assumes second-countable S and subbase-in S U
  shows borel-of S = sigma (topspace S) U
  unfolding borel-of-def
proof(rule sigma-eqI)
  show {U. openin S U}  $\subseteq$  Pow (topspace S)
  by (simp add: subset-Pow-Union topspace-def)
next
  show  $U \subseteq$  Pow (topspace S)
  using subbase-in-subset[OF assms(2)] by auto
next
  interpret s: sigma-algebra topspace S sigma-sets (topspace S) U
  using subbase-in-subset[OF assms(2)] by(auto intro!: sigma-algebra-sigma-sets)
  obtain O where ho: countable O base-in S O
  using assms(1) by(auto simp: second-countable-base-in)
  show sigma-sets (topspace S) {U. openin S U} = sigma-sets (topspace S) U
  proof(rule sigma-sets-eqI)
  fix U
  assume  $U \in \{U. \text{openin } S U\}$ 
  then have generate-topology-on U U
  using assms(2) by(simp add: subbase-in-def openin-topology-generated-by-iff)
  thus  $U \in \text{sigma-sets (topspace S) } U$ 
  proof induction
  case (UN K)
  with ho(2) obtain V where hv:
     $\bigwedge k. k \in K \implies V k \subseteq O \wedge k. k \in K \implies \bigcup (V k) = k$ 
  by(simp add: base-in-def openin-topology-generated-by-iff[symmetric] assms(2)[simplified
subbase-in-def,symmetric]) metis
  define Uk where  $Uk = (\bigcup_{k \in K}. V k)$ 
  have 0:countable Uk
  using hv by(auto intro!: countable-subset[OF - ho(1)] simp: Uk-def)
  have  $\bigcup Uk = (\bigcup_{A \in Uk}. A)$  by auto
  also have ... =  $\bigcup K$ 

```

```

    unfolding  $\mathcal{U}k$ -def UN-simps by (simp add: hv(2))
    finally have  $1: \bigcup \mathcal{U}k = \bigcup K$  .
    have  $\forall b \in \mathcal{U}k. \exists k \in K. b \subseteq k$ 
    using hv by (auto simp:  $\mathcal{U}k$ -def)
    then obtain  $V'$  where  $hv': \bigwedge b. b \in \mathcal{U}k \implies V' b \in K$  and  $\bigwedge b. b \in \mathcal{U}k \implies$ 
 $b \subseteq V' b$ 
    by metis
    then have  $(\bigcup b \in \mathcal{U}k. V' b) \subseteq \bigcup K \cup \mathcal{U}k \subseteq (\bigcup b \in \mathcal{U}k. V' b)$ 
    by auto
    then have  $\bigcup K = (\bigcup b \in \mathcal{U}k. V' b)$ 
    unfolding 1 by auto
    also have  $\dots \in \text{sigma-sets (topspace } S) \mathcal{U}$ 
    using  $hv'$  UN by (auto intro!: s.countable-UN' simp: 0)
    finally show  $\bigcup K \in \text{sigma-sets (topspace } S) \mathcal{U}$  .
  qed auto
next
fix  $U$ 
assume  $U \in \mathcal{U}$ 
from  $\text{assms}(2)[\text{simplified subbase-in-def}] \text{openin-topology-generated-by-iff generate-topology-on.Basis[OF this]}$ 
show  $U \in \text{sigma-sets (topspace } S) \{U. \text{openin } S U\}$ 
by auto
qed
qed

```

Abstract topology version $\text{borel} \otimes_M \text{borel} = \text{borel}$.

lemma *borel-of-prod*:

```

  assumes second-countable  $S$  and second-countable  $S'$ 
  shows borel-of  $S \otimes_M \text{borel-of } S' = \text{borel-of (prod-topology } S S')$ 
proof -
  have borel-of  $S \otimes_M \text{borel-of } S' = \text{sigma (topspace } S \times \text{topspace } S') \{a \times b \mid a$ 
 $b. a \in \{a. \text{openin } S a\} \wedge b \in \{b. \text{openin } S' b\}\}$ 
  proof -
    obtain  $\mathcal{O} \mathcal{O}'$  where ho:
    countable  $\mathcal{O}$  base-in  $S$   $\mathcal{O}$  countable  $\mathcal{O}'$  base-in  $S'$   $\mathcal{O}'$ 
    using assms by (auto simp: second-countable-base-in)
    show ?thesis
    unfolding borel-of-def
    apply (rule sigma-prod)
    using topology-generated-by-topspace[of  $\mathcal{O}$ , simplified base-is-subbase[OF ho(2), simplified
subbase-in-def, symmetric]] topology-generated-by-topspace[of  $\mathcal{O}'$ , simplified base-is-subbase[OF
 $ho(4)$ , simplified subbase-in-def, symmetric]]
    base-in-openin[OF ho(2)] base-in-openin[OF  $ho(4)$ ]
    by (auto intro!:  $exI[\text{where } x=\mathcal{O}] exI[\text{where } x=\mathcal{O}']$  simp: ho subset-Pow-Union
topspace-def)
  qed
  also have  $\dots = \text{borel-of (prod-topology } S S')$ 
  using borel-of-second-countable'[OF prod-topology-second-countable[OF assms], simplified
subbase-in-def, OF prod-topology-generated-by-open]

```

by simp
 finally show ?thesis .
 qed

lemma product-borel-of-measurable:

assumes $i \in I$
 shows $(\lambda x. x i) \in (\text{borel-of } (\text{product-topology } S I)) \rightarrow_M \text{borel-of } (S i)$
 by(auto intro!: continuous-map-measurable simp: assms)

Abstract topology version of sets $(\text{Pi}_M \text{ UNIV } (\lambda-. \text{borel})) \subseteq \text{sets borel}$

lemma sets-PiM-subset-borel-of:

sets $(\text{Pi}_M i \in I. \text{borel-of } (S i)) \subseteq \text{sets } (\text{borel-of } (\text{product-topology } S I))$

proof –

have *: $(\text{Pi}_E i \in I. X i) \in \text{sets } (\text{borel-of } (\text{product-topology } S I))$ if [measurable]: $\bigwedge i. X i \in \text{sets } (\text{borel-of } (S i))$ finite $\{i. X i \neq \text{topspace } (S i)\}$ for X

proof –

note [measurable] = product-borel-of-measurable

define I' where $I' = \{i. X i \neq \text{topspace } (S i)\} \cap I$

have finite I' unfolding I' -def using that by simp

have $(\text{Pi}_E i \in I. X i) = (\bigcap i \in I'. (\lambda x. x i) -' (X i) \cap \text{space } (\text{borel-of } (\text{product-topology } S I))) \cap \text{space } (\text{borel-of } (\text{product-topology } S I))$

proof(standard;standard)

fix x

assume $x \in \text{Pi}_E I X$

then show $x \in (\bigcap i \in I'. (\lambda x. x i) -' X i \cap \text{space } (\text{borel-of } (\text{product-topology } S I))) \cap \text{space } (\text{borel-of } (\text{product-topology } S I))$

using sets.sets-into-space[OF that(1)] by(auto simp: PiE-def I'-def Pi-def space-borel-of)

next

fix x

assume $1: x \in (\bigcap i \in I'. (\lambda x. x i) -' X i \cap \text{space } (\text{borel-of } (\text{product-topology } S I))) \cap \text{space } (\text{borel-of } (\text{product-topology } S I))$

have $x i \in X i$ if $hi: i \in I$ for i

proof –

consider $i \in I' \wedge I' \neq \{\} \mid i \notin I' \wedge I' = \{\} \mid i \notin I' \wedge I' \neq \{\}$ by auto

then show ?thesis

apply cases

using sets.sets-into-space[OF $\langle \bigwedge i. X i \in \text{sets } (\text{borel-of } (S i)) \rangle$] 1 that

by(auto simp: space-borel-of I'-def)

qed

then show $x \in \text{Pi}_E I X$

using 1 by(auto simp: space-borel-of)

qed

also have $\dots \in \text{sets } (\text{borel-of } (\text{product-topology } S I))$

using that $\langle \text{finite } I' \rangle$ by(auto simp: I'-def)

finally show ?thesis .

qed

then have $\{\text{Pi}_E I X \mid X. (\forall i. X i \in \text{sets } (\text{borel-of } (S i))) \wedge \text{finite } \{i. X i \neq \text{space } (\text{borel-of } (S i))\}\} \subseteq \text{sets } (\text{borel-of } (\text{product-topology } S I))$

by(*auto simp: space-borel-of*)
show *?thesis unfolding sets-PiM-finite*
by(*rule sets.sigma-sets-subset',fact*) (*simp add: borel-of-open[OF openin-topospace,*
of product-topology S I,simplified] space-borel-of)
qed

Abstract topology version of *sets (Pi_M UNIV (λi. borel)) = sets borel*.

lemma *sets-PiM-equal-borel-of*:

assumes *countable I and* $\bigwedge i. i \in I \implies \text{second-countable } (S\ i)$

shows *sets (Π_M i∈I. borel-of (S i)) = sets (borel-of (product-topology S I))*

proof

obtain *K where hk:*

countable K base-in (product-topology S I) K

$\bigwedge k. k \in K \implies \exists X. (k = (\prod_{E} i \in I. X\ i)) \wedge (\forall i. \text{openin } (S\ i) (X\ i)) \wedge \text{finite } \{i. X\ i \neq \text{topspace } (S\ i)\} \wedge \{i. X\ i \neq \text{topspace } (S\ i)\} \subseteq I$

using *product-topology-countable-base-in[OF assms(1)] assms(2)*

by force

have **:k ∈ sets (Π_M i∈I. borel-of (S i)) if k ∈ K for k*

proof –

obtain *X where H: k = (Π_E i∈I. X i) ∧ i. openin (S i) (X i) finite {i. X i ≠ topspace (S i)} {i. X i ≠ topspace (S i)} ⊆ I*

using *hk(3)[OF ‹k ∈ K›] by blast*

show *?thesis unfolding H(1) sets-PiM-finite*

using *borel-of-open[OF H(2)] H(3) by(auto simp: space-borel-of)*

qed

have *** : U ∈ sets (Π_M i∈I. borel-of (S i)) if openin (product-topology S I) U*
for *U*

proof –

obtain *B where B ⊆ K U = (∪ B)*

using *‹openin (product-topology S I) U› ‹base-in (product-topology S I) K›*

by (*metis base-in-def*)

have *countable B using ‹B ⊆ K› ‹countable K› countable-subset by blast*

moreover have *k ∈ sets (Π_M i∈I. borel-of (S i)) if k ∈ B for k*

using *‹B ⊆ K› * that by auto*

ultimately show *?thesis unfolding ‹U = (∪ B)› by auto*

qed

have *sigma-sets (topspace (product-topology S I)) {U. openin (product-topology S I) U} ⊆ sets (Π_M i∈I. borel-of (S i))*

apply (*rule sets.sigma-sets-subset'*) **using** *** by(auto intro!: sets-PiM-I-countable[OF assms(1)] simp: borel-of-open[OF openin-topospace])*

thus *sets (borel-of (product-topology S I)) ⊆ sets (Π_M i∈I. borel-of (S i))*

by (*simp add: subset-Pow-Union topspace-def borel-of-def*)

qed(*rule sets-PiM-subset-borel-of*)

lemma *homeomorphic-map-borel-isomorphic*:

assumes *homeomorphic-map X Y f*

shows *measurable-isomorphic-map (borel-of X) (borel-of Y) f*

proof –

obtain *g where homeomorphic-maps X Y f g*

using *assms* **by**(*auto simp: homeomorphic-map-maps*)
hence *continuous-map X Y f continuous-map Y X g*
 $\bigwedge x. x \in \text{topspace } X \implies g (f x) = x$
 $\bigwedge y. y \in \text{topspace } Y \implies f (g y) = y$
by(*auto simp: homeomorphic-maps-def*)
thus *?thesis*
by(*auto intro!: measurable-isomorphic-map-byWitness dest: continuous-map-measurable simp: space-borel-of*)
qed

lemma *homeomorphic-space-measurable-isomorphic:*
assumes *S homeomorphic-space T*
shows *borel-of S measurable-isomorphic borel-of T*
using *homeomorphic-map-borel-isomorphic[of S T] assms* **by**(*auto simp: measurable-isomorphic-def homeomorphic-space*)

lemma *measurable-isomorphic-borel-map:*
assumes *sets M = sets (borel-of S) and f: measurable-isomorphic-map M N f*
shows $\exists S'. \text{homeomorphic-map } S S' f \wedge \text{sets } N = \text{sets (borel-of } S')$
proof –

obtain *g* **where** $fg: f \in M \rightarrow_M N \ g \in N \rightarrow_M M \ \bigwedge x. x \in \text{space } M \implies g (f x) = x$
 $\bigwedge y. y \in \text{space } N \implies f (g y) = y \ \bigwedge A. A \in \text{sets } M \implies f^{-1} A \in \text{sets } N \ \bigwedge A. A \in \text{sets } N \implies g^{-1} A \in \text{sets } M$ *bij-betw g (space N) (space M)*

using *measurable-isomorphic-mapD'[OF f]* **by** *metis*

have *g: measurable-isomorphic-map N M g*

by(*auto intro!: measurable-isomorphic-map-byWitness fg*)

have *g': bij-betw g (space N) (topspace S)*

using *fg(7) sets-eq-imp-space-eq[OF assms(1)]* **by**(*auto simp: space-borel-of*)

show *?thesis*

proof(*intro exI[where x=pullback-topology (space N) g S] conjI*)

have [*simp*]: $\{U. \text{openin (pullback-topology (space N) g S) } U\} = (\cdot)^{-1} f^{-1} \{U. \text{openin } S \ U\}$

unfolding *openin-pullback-topology'[OF g']*

proof *safe*

fix *u*

assume *u: openin S u*

then have $1: u \subseteq \text{space } M$

by(*simp add: sets-eq-imp-space-eq[OF assms(1)] space-borel-of openin-subset*)

with *fg(3)* **have** $g^{-1} f^{-1} u = u$

by(*fastforce simp: image-def*)

with *u* **show** *openin S (g^{-1} f^{-1} u)* **by** *simp*

fix *x*

assume $x \in u$

with $1 \ fg(1)$ **show** $f x \in \text{space } N$ **by**(*auto simp: measurable-space*)

next

fix *u*

assume *openin S (g^{-1} u) u \subseteq \text{space } N*

with *fg(4)* **show** $u \in (\cdot)^{-1} f^{-1} \{U. \text{openin } S \ U\}$

by(*auto simp: image-def intro!: exI[where x=g^{-1} u] (metis in-mono)*)

```

qed
have [simp]:g -' topspace S ∩ space N = space N
  using bij-betw-imp-surj-on g' by blast
show sets N = sets (borel-of (pullback-topology (space N) g S))
  by(auto simp: sets-borel-of topspace-pullback-topology intro!: measurable-isomorphic-map-sigma-sets[OF
assms(1)][simplified sets-borel-of space-borel-of[symmetric] sets-eq-imp-space-eq[OF
assms(1),symmetric]] f)
next
show homeomorphic-map S (pullback-topology (space N) g S) f
  proof(rule homeomorphic-maps-imp-map[where g=g])
    obtain f' where f':homeomorphic-maps (pullback-topology (space N) g S) S
g f'
      using topology-from-bij(1)[OF g'] homeomorphic-map-maps by blast
    have f'2:f' y = f y if y:y ∈ topspace S for y
    proof -
      have [simp]:g -' topspace S ∩ space N = space N
        using bij-betw-imp-surj-on g' by blast
      obtain x where x ∈ space N y = g x
        using g' y by(auto simp: bij-betw-def image-def)
      thus ?thesis
    using fg(4) f' by(auto simp: homeomorphic-maps-def topspace-pullback-topology)
    qed
  thus homeomorphic-maps S (pullback-topology (space N) g S) f g
  by(auto intro!: homeomorphic-maps-eq[OF f'] simp: homeomorphic-maps-sym[of
S])
qed
qed
qed

```

```

lemma measurable-isomorphic-borels:
  assumes sets M = sets (borel-of S) M measurable-isomorphic N
  shows ∃ S'. S homeomorphic-space S' ∧ sets N = sets (borel-of S')
  using measurable-isomorphic-borel-map[OF assms(1)] assms(2) homeomorphic-map-maps
  by(fastforce simp: measurable-isomorphic-def homeomorphic-space-def )

```

end

1.3 Lemmas for Abstract Metric Spaces

```

theory Set-Based-Metric-Space
  imports Lemmas-StandardBorel
begin

```

We prove additional lemmas related to set-based metric spaces.

1.3.1 Basic Lemmas

```

lemma
  assumes Metric-space M d ∧ x y. x ∈ M ⇒ y ∈ M ⇒ d x y = d' x y

```



```

    and  $\bigwedge x y. d' x y = d' y x \wedge x y. d' x y \geq 0$ 
  shows Metric-space-eq: Metric-space  $M d'$ 
  and Metric-space-eq-mtopology: Metric-space.mtopology  $M d = \text{Metric-space.mtopology}$ 
   $M d'$ 
    and Metric-space-eq-mcomplete: Metric-space.mcomplete  $M d \longleftrightarrow \text{Metric-}$ 
space.mcomplete  $M d'$ 
  proof –
    interpret  $m1$ : Metric-space  $M d$  by fact
    show Metric-space  $M d'$ 
      using assms by(auto simp: Metric-space-def)
    then interpret  $m2$ : Metric-space  $M d'$  .
    have [simp]:  $m1.mball\ x\ e = m2.mball\ x\ e$  for  $x\ e$ 
      using assms by(auto simp:  $m1.mball-def\ m2.mball-def$ )
    show  $1:m1.mtopology = m2.mtopology$ 
      by(auto simp: topology-eq  $m1.openin-mtopology\ m2.openin-mtopology$ )
    show  $m1.mcomplete = m2.mcomplete$ 
      by(auto simp:  $1\ m1.mcomplete-def\ m2.mcomplete-def\ m1.MCauchy-def\ m2.MCauchy-def$ 
assms(2) in-mono)
  qed

```

```

context Metric-space
begin

```

```

lemma mtopology-base-in-balls: base-in mtopology  $\{mball\ a\ \varepsilon \mid a\ \varepsilon. a \in M \wedge \varepsilon > 0\}$ 
}
proof –
  have  $1:\bigwedge x. x \in \{mball\ a\ \varepsilon \mid a\ \varepsilon. a \in M \wedge \varepsilon > 0\} \implies \text{openin}\ mtopology\ x$ 
    by auto
  show ?thesis
    unfolding base-in-def2[of  $\{mball\ a\ \varepsilon \mid a\ \varepsilon. a \in M \wedge \varepsilon > 0\}$ , OF 1, simplified]
    by (metis centre-in-mball-iff in-mono openin-mtopology)
  qed

```

```

lemma closedin-metric2: closedin mtopology  $C \longleftrightarrow C \subseteq M \wedge (\forall x. x \in C \longleftrightarrow$ 
 $(\forall \varepsilon > 0. mball\ x\ \varepsilon \cap C \neq \{\}))$ 

```

```

proof
  assume  $h$ : closedin mtopology  $C$ 
  have  $1: C \subseteq M$ 
    using Metric-space.closedin-metric Metric-space-axioms h by blast
  show  $C \subseteq M \wedge (\forall x. x \in C \longleftrightarrow (\forall \varepsilon > 0. mball\ x\ \varepsilon \cap C \neq \{\}))$ 
  proof safe
    fix  $\varepsilon\ x$ 
    assume  $x \in C\ (0 :: real) < \varepsilon\ mball\ x\ \varepsilon \cap C = \{\}$ 
    with  $1$  show False
      by blast
  next
    fix  $x$ 
    assume  $\forall \varepsilon > 0. mball\ x\ \varepsilon \cap C \neq \{\}$ 
    hence  $\exists xn. xn \in mball\ x\ (1 / real\ (Suc\ n)) \cap C$  for  $n$ 

```

```

by (meson all-not-in-conv divide-pos-pos of-nat-0-less-iff zero-less-Suc zero-less-one)
then obtain xn where xn:  $\bigwedge n. xn\ n \in \text{mball } x\ (1 / \text{real } (\text{Suc } n)) \cap C$ 
  by metis
hence xxn:  $x \in M\ \text{range } xn \subseteq C$ 
  using xn by auto
have limitin mtopology xn x sequentially
  unfolding limitin-metric eventually-sequentially
proof safe
  fix  $\varepsilon$ 
  assume  $(0 :: \text{real}) < \varepsilon$ 
  then obtain N where hN:  $1 / \text{real } (\text{Suc } N) < \varepsilon$ 
    using nat-approx-posE by blast
  show  $\exists N. \forall n \geq N. xn\ n \in M \wedge d(xn\ n)\ x < \varepsilon$ 
  proof (safe intro!: exI[where x=N])
    fix n
    assume n[arith]:  $N \leq n$ 
    then have  $1 / \text{real } (\text{Suc } n) < \varepsilon$ 
    by (metis Suc-le-mono hN inverse-of-nat-le nat.distinct(1) order-le-less-trans)
    with xn[of n] show  $d(xn\ n)\ x < \varepsilon$ 
    by (simp add: commute)
  qed (use xxn 1 in auto)
qed fact
with h 1 xxn show  $x \in C$ 
  by (auto simp: metric-closedin-iff-sequentially-closed)
qed (use 1 in auto)
next
assume  $C \subseteq M \wedge (\forall x. (x \in C) \longleftrightarrow (\forall \varepsilon > 0. \text{mball } x\ \varepsilon \cap C \neq \{\}))$ 
hence h:  $C \subseteq M \wedge x. (x \in C) \longleftrightarrow (\forall \varepsilon > 0. \text{mball } x\ \varepsilon \cap C \neq \{\})$ 
  by simp-all
show closedin mtopology C
  unfolding metric-closedin-iff-sequentially-closed
proof safe
  fix xn x
  assume h':  $\text{range } xn \subseteq C\ \text{limitin mtopology xn x sequentially}$ 
  hence  $x \in M$  by (simp add: limitin-mspace)
  have  $\text{mball } x\ \varepsilon \cap C \neq \{\}$  if  $\varepsilon > 0$  for  $\varepsilon$ 
  proof -
    obtain N where hN:  $\bigwedge n. n \geq N \implies d(xn\ n)\ x < \varepsilon$ 
      using h'(2)  $\langle \varepsilon > 0 \rangle$  limit-metric-sequentially by blast
    have  $xn\ N \in \text{mball } x\ \varepsilon \cap C$ 
      using h'(1) hN[of N]  $\langle x \in M \rangle$  commute h(1) by fastforce
    thus  $\text{mball } x\ \varepsilon \cap C \neq \{\}$  by auto
  qed
  with h(2)[of x] show  $x \in C$  by simp
qed (use h(1) in auto)
qed

```

lemma openin-mtopology2:

$\text{openin mtopology } U \longleftrightarrow U \subseteq M \wedge (\forall xn\ x. \text{limitin mtopology xn x sequentially} \wedge$

```

x ∈ U → (∃ N. ∀ n ≥ N. xn n ∈ U))
  unfolding openin-mtopology
proof safe
  fix xn x
  assume h: ∀ x. x ∈ U → (∃ r > 0. mball x r ⊆ U) limitin mtopology xn x
  sequentially x ∈ U U ⊆ M
  then obtain r where r: r > 0 mball x r ⊆ U
  by auto
  with h(2) obtain N where N: ∧ n. n ≥ N ⇒ xn n ∈ M ∧ n. n ≥ N ⇒ d
  (xn n) x < r
  by (metis limit-metric-sequentially)
  with h have ∃ N. ∀ n ≥ N. xn n ∈ mball x r
  by (auto intro!: exI[where x=N] simp: commute)
  with r show ∃ N. ∀ n ≥ N. xn n ∈ U
  by blast
next
fix x
assume h: U ⊆ M ∀ xn x. limitin mtopology xn x sequentially ∧ x ∈ U → (∃ N.
∀ n ≥ N. xn n ∈ U) x ∈ U
show ∃ r > 0. mball x r ⊆ U
proof (rule ccontr)
  assume ¬ (∃ r > 0. mball x r ⊆ U)
  then have ∀ n. ∃ xn ∈ mball x (1 / Suc n). xn ∉ U
  by (meson of-nat-0-less-iff subsetI zero-less-Suc zero-less-divide-1-iff)
  then obtain xn where xn: ∧ n. xn n ∈ mball x (1 / Suc n) ∧ n. xn n ∉ U
  by metis
  have limitin mtopology xn x sequentially
  unfolding limit-metric-sequentially
proof safe
  fix e :: real
  assume e: 0 < e
  then obtain N where N: 1 / real (Suc N) < e
  using nat-approx-posE by blast
  show ∃ N. ∀ n ≥ N. xn n ∈ M ∧ d (xn n) x < e
  proof (safe intro!: exI[where x=N])
    fix n
    assume n: n ≥ N
    then have 1 / Suc n < e
    by (metis N Suc-le-mono inverse-of-nat-le nat.distinct(1) order-le-less-trans)
    thus d (xn n) x < e
    using xn(1)[of n] by (auto simp: commute)
  qed (use xn in auto)
  qed (use h in auto)
  with h(2,3) xn(2) show False
  by auto
qed
qed

```

lemma closure-of-mball: mtopology closure-of mball a e ⊆ mcball a e

by (simp add: closure-of-minimal mball-subset-mcball)

lemma interior-of-mcball: $mball\ a\ e \subseteq mtopology\ interior\ of\ mcball\ a\ e$
 by (simp add: interior-of-maximal-eq mball-subset-mcball)

lemma isolated-points-of-mtopology:
 $mtopology\ isolated\ points\ of\ A = \{x \in M \cap A. \forall xn. range\ xn \subseteq A \wedge limitin\ mtopology\ xn\ x\ sequentially \longrightarrow (\exists no. \forall n \geq no. xn\ n = x)\}$

proof safe
 fix $x\ xn$
 assume $h: x \in mtopology\ isolated\ points\ of\ A\ limitin\ mtopology\ xn\ x\ sequentially$
 $range\ xn \subseteq A$
 then have $ha: x \in topspace\ mtopology\ x \in A \exists U. x \in U \wedge openin\ mtopology\ U$
 $\wedge U \cap (A - \{x\}) = \{\}$
 by (simp-all add: in-isolated-points-of)
 then obtain U where $u: x \in U\ openin\ mtopology\ U\ U \cap (A - \{x\}) = \{\}$
 by auto
 then obtain e where $e: e > 0\ mball\ x\ e \subseteq U$
 by (meson openin-mtopology)
 then obtain N where $\bigwedge n. n \geq N \implies xn\ n \in mball\ x\ e$
 using $h(2)$ commute limit-metric-sequentially by fastforce
 thus $\exists no. \forall n \geq no. xn\ n = x$
 using $h(3)\ e(2)\ u(3)$ by (fastforce intro!: exI[where $x=N$])
 qed (auto simp: derived-set-of-sequentially isolated-points-of-def, blast)

lemma perfect-set-mball-infinite:
 assumes perfect-set mtopology $A\ a \in A\ e > 0$
 shows infinite (mball $a\ e$)

proof safe
 assume $h: finite\ (mball\ a\ e)$
 have $a \in M$
 using $assms\ perfect\ setD(2)[OF\ assms(1)]$ by auto
 have $\exists e > 0. mball\ a\ e = \{a\}$
proof –
 consider $mball\ a\ e = \{a\} \mid \{a\} \subset mball\ a\ e$
 using $\langle a \in M \rangle\ assms(3)$ by blast
 thus ?thesis
proof cases
 case 1
 with $assms$ show ?thesis by auto
 next
 case 2
 then have $nen: \{d\ a\ b \mid b. b \in mball\ a\ e \wedge a \neq b\} \neq \{\}$
 by auto
 have $fin: finite\ \{d\ a\ b \mid b. b \in mball\ a\ e \wedge a \neq b\}$
 using h by (auto simp del: in-mball)
 define e' where $e' \equiv Min\ \{d\ a\ b \mid b. b \in mball\ a\ e \wedge a \neq b\}$
 have $e' > 0$
 using $mdist\ pos\ eq[OF\ \langle a \in M \rangle]$ by (simp add: e'-def Min-gr-iff[OF fin nen])

```

del: in-mball) auto
  have bd:  $\bigwedge b. b \in \text{mball } a \ e \implies a \neq b \implies e' \leq d \ a \ b$ 
    by (auto simp: e'-def Min-le-iff[OF fin nen] simp del: in-mball)
  have e'  $\leq e$ 
    unfolding e'-def Min-le-iff[OF fin nen]
    using nen by auto
  show ?thesis
  proof (safe intro!: exI[where x=e'])
    fix x
    assume x:  $x \in \text{mball } a \ e'$ 
    then show  $x = a$ 
      using  $\langle e' \leq e \rangle$  bd by fastforce
    qed (use  $\langle a \in M \rangle \langle e' > 0 \rangle$  in auto)
  qed
  then obtain e' where e':  $e' > 0 \ \text{mball } a \ e' = \{a\}$  by auto
  show False
    using perfect-setD(3)[OF assms(1,2) centre-in-mball-iff[of a e', THEN iffD2]]
     $\langle a \in M \rangle \langle e' > 0 \rangle e'(\mathcal{Q})$  by blast
  qed

```

lemma *MCauchy-dist-Cauchy*:

```

assumes MCauchy  $xn \ \text{MCauchy } yn$ 
shows Cauchy  $(\lambda n. d \ (xn \ n) \ (yn \ n))$ 
unfolding metric-space-class.Cauchy-altdef2 dist-real-def
proof safe
  have h:  $\bigwedge n. xn \ n \in M \ \bigwedge n. yn \ n \in M$ 
    using assms by (auto simp: MCauchy-def)
  fix e :: real
  assume e:  $0 < e$ 
  with assms obtain  $N1 \ N2$  where  $N: \bigwedge n \ m. n \geq N1 \implies m \geq N1 \implies d \ (xn \ n) \ (xn \ m) < e / 2 \ \bigwedge n \ m. n \geq N2 \implies m \geq N2 \implies d \ (yn \ n) \ (yn \ m) < e / 2$ 
    by (metis MCauchy-def zero-less-divide-iff zero-less-numeral)
  define  $N$  where  $N \equiv \max \ N1 \ N2$ 
  then have  $N': N \geq N1 \ N \geq N2$  by auto
  show  $\exists N. \forall n \geq N. |d \ (xn \ n) \ (yn \ n) - d \ (xn \ N) \ (yn \ N)| < e$ 
  proof (safe intro!: exI[where x=N])
    fix n
    assume n:  $N \leq n$ 
    have  $d \ (xn \ n) \ (yn \ n) \leq d \ (xn \ n) \ (xn \ N) + d \ (xn \ N) \ (yn \ N) + d \ (yn \ N) \ (yn \ n)$ 
       $d \ (xn \ N) \ (yn \ N) \leq d \ (xn \ N) \ (xn \ n) + d \ (xn \ n) \ (yn \ n) + d \ (yn \ n) \ (yn \ N)$ 
      using triangle[OF h(1)[of n] h(1)[of N] h(2)[of n]] triangle[OF h(1)[of N] h(2)[of N] h(2)[of n]]
      triangle[OF h(1)[of N] h(2)[of n] h(2)[of N]] triangle[OF h(1)[of N] h(1)[of n] h(2)[of n]] by auto
    thus  $|d \ (xn \ n) \ (yn \ n) - d \ (xn \ N) \ (yn \ N)| < e$ 
      using  $N(1)$ [OF  $N'(1)$  order.trans[OF  $N'(1)$  n]]  $N(2)$ [OF  $N'(2)$  order.trans[OF  $N'(2)$  n]]  $N(1)$ [OF order.trans[OF  $N'(1)$  n]  $N'(1)$ ]  $N(2)$ [OF order.trans[OF  $N'(2)$  n]  $N'(2)$ ]

```

```

n] N'(2)]
  by auto
  qed
qed

```

1.3.2 Dense in Metric Spaces

abbreviation $mdense \equiv dense\text{-in } mtopology$

<https://people.bath.ac.uk/mw2319/ma30252/sec-dense.html>

lemma $mdense\text{-def}$:

$mdense\ U \longleftrightarrow U \subseteq M \wedge (\forall x \in M. \forall \varepsilon > 0. mball\ x\ \varepsilon \cap U \neq \{\})$

proof *safe*

assume $h: U \subseteq M (\forall x \in M. \forall \varepsilon > 0. mball\ x\ \varepsilon \cap U \neq \{\})$

show $dense\text{-in } mtopology\ U$

proof(*rule dense-inI*)

fix V

assume $h': openin\ mtopology\ V\ V \neq \{\}$

then obtain x **where** $1: x \in V$ **by** *auto*

then obtain ε **where** $2: \varepsilon > 0\ mball\ x\ \varepsilon \subseteq V$

by (*meson h'(1) openin-mtopology*)

have $mball\ x\ \varepsilon \cap U \neq \{\}$

using $h\ 1\ 2\ openin\text{-subset}[OF\ h'(1)]$

by (*auto simp del: in-mball*)

thus $U \cap V \neq \{\}$ **using** 2 **by** *auto*

qed(*use h in auto*)

next

fix $x\ \varepsilon$

assume $h: x \in M\ (0 :: real) < \varepsilon\ mball\ x\ \varepsilon \cap U = \{\}\ mdense\ U$

then have $mball\ x\ \varepsilon \neq \{\}\ openin\ mtopology\ (mball\ x\ \varepsilon)$

by *auto*

with $h(4)$ **have** $mball\ x\ \varepsilon \cap U \neq \{\}$

by(*auto simp: dense-in-def*)

with $h(3)$ **show** *False*

by *simp*

qed(*auto simp: dense-in-def*)

corollary $mdense\text{-balls-cover}$:

assumes $mdense\ U$ **and** $e > 0$

shows $(\bigcup u \in U. mball\ u\ e) = M$

using *assms[simplified mdense-def] commute* **by** *fastforce*

lemma $mdense\text{-empty-iff}$: $mdense\ \{\} \longleftrightarrow M = \{\}$

by(*auto simp: mdense-def*) (*use zero-less-one in blast*)

lemma $mdense\text{-M}$: $mdense\ M$

by(*auto simp: mdense-def*)

lemma $mdense\text{-def2}$:

```

mdense U  $\longleftrightarrow$  U  $\subseteq$  M  $\wedge$  ( $\forall x \in M. \forall \varepsilon > 0. \exists y \in U. d x y < \varepsilon$ )
proof safe
  fix x e
  assume h: mdense U and hxe: x  $\in$  M (0 :: real) < e
  then have x  $\in$  ( $\bigcup u \in U. mball u e$ )
    by(simp add: mdense-balls-cover)
  thus  $\exists y \in U. d x y < e$ 
    by (fastforce simp: commute)
qed(fastforce simp: mdense-def)+

lemma mdense-def3:
  mdense U  $\longleftrightarrow$  U  $\subseteq$  M  $\wedge$  ( $\forall x \in M. \exists xn. range xn \subseteq U \wedge limitin mtopology xn x$ 
  sequentially)
  unfolding mdense-def
proof safe
  fix x
  assume h: U  $\subseteq$  M  $\forall x \in M. \forall \varepsilon > 0. mball x \varepsilon \cap U \neq \{\}$  x  $\in$  M
  then have  $\bigwedge n. mball x (1 / (real n + 1)) \cap U \neq \{\}$ 
    by auto
  hence  $\forall n. \exists k. k \in mball x (1 / (real n + 1)) \cap U$  by (auto simp del: in-mball)
  hence  $\exists a. \forall n. a n \in mball x (1 / (real n + 1)) \cap U$  by(rule choice)
  then obtain xn where xn:  $\bigwedge n. xn n \in mball x (1 / (real n + 1)) \cap U$ 
    by auto
  show  $\exists xn. range xn \subseteq U \wedge limitin mtopology xn x$  sequentially
    unfolding limitin-metric eventually-sequentially
  proof(safe intro!: exI[where x=xn])
    fix  $\varepsilon :: real$ 
    assume he:0 <  $\varepsilon$ 
    then obtain N where hn: 1 /  $\varepsilon < real N$ 
      using reals-Archimedean2 by blast
    have hn': 0 < real N
      by(rule ccontr) (use hn he in fastforce)
    hence 1 / real N <  $\varepsilon$ 
      using he hn by (metis divide-less-eq mult.commute)
    hence hn'':1 / (real N + 1) <  $\varepsilon$ 
      using hn' by(auto intro!: order.strict-trans[OF linordered-field-class.divide-strict-left-mono[of
  real N real N + 1 1]])
    show  $\exists N. \forall n \geq N. xn n \in M \wedge d (xn n) x < \varepsilon$ 
      proof(safe intro!: exI[where x=N])
        fix n
        assume N  $\leq$  n
        then have 1:1 / (real n + 1)  $\leq$  1 / (real N + 1)
          using hn' by(auto intro!: linordered-field-class.divide-left-mono)
        show d (xn n) x <  $\varepsilon$ 
          using xn[of n] order.strict-trans1[OF 1 hn''] by (auto simp: commute)
        qed(use xn in auto)
      qed(use xn h in auto)
  next
  fix x and e :: real

```

assume $h: U \subseteq M \ \forall x \in M. \exists xn. \text{range } xn \subseteq U \wedge \text{limitin } mtopology \ xn \ x \text{ sequentially}$
 $x \in M \ 0 < e \ \text{mball } x \ e \cap U = \{\}$
then obtain xn **where** $xn: \text{range } xn \subseteq U \ \text{limitin } mtopology \ xn \ x \text{ sequentially}$
by *auto*
with $h(4)$ **obtain** N **where** $N: \bigwedge n. n \geq N \implies d(xn \ n) \ x < e$
by (*meson limit-metric-sequentially*)
have $xn \ N \in \text{mball } x \ e \cap U$
using $N[\text{of } N] \ xn(1) \ h(1,3)$ **by** (*auto simp: commute*)
with $h(5)$ **show** *False* **by** *simp*
qed

Diameter

definition $mdiameter :: 'a \text{ set} \Rightarrow \text{ennreal}$ **where**
 $mdiameter \ A \equiv \bigsqcup \{ \text{ennreal } (d \ x \ y) \mid x \ y. \ x \in A \cap M \wedge y \in A \cap M \}$

lemma *mdiameter-empty[simp]*:
 $mdiameter \ \{\} = 0$
by (*simp add: mdiameter-def bot-ennreal*)

lemma *mdiameter-def2*:
assumes $A \subseteq M$
shows $mdiameter \ A = \bigsqcup \{ \text{ennreal } (d \ x \ y) \mid x \ y. \ x \in A \wedge y \in A \}$
using *assms* **by** (*auto simp: mdiameter-def (meson subset-eq)*)

lemma *mdiameter-subset*:
assumes $A \subseteq B$
shows $mdiameter \ A \leq mdiameter \ B$
unfolding *mdiameter-def* **using** *assms* **by** (*auto intro!: Sup-subset-mono*)

lemma *mdiameter-cball-leq*: $mdiameter \ (\text{mcball } a \ \varepsilon) \leq \text{ennreal } (2 * \varepsilon)$
unfolding *Sup-le-iff mdiameter-def*
proof *safe*

fix $x \ y$
assume $h: x \in \text{mcball } a \ \varepsilon \ y \in \text{mcball } a \ \varepsilon \ x \in M \ y \in M$
have $d \ x \ y \leq 2 * \varepsilon$
using $h(1) \ h(2)$ *triangle''* **by** *fastforce*
thus $\text{ennreal } (d \ x \ y) \leq \text{ennreal } (2 * \varepsilon)$
using *ennreal-leI* **by** *blast*
qed

lemma *mdiameter-ball-leq*:
 $mdiameter \ (\text{mball } a \ \varepsilon) \leq \text{ennreal } (2 * \varepsilon)$
using *mdiameter-subset[OF mball-subset-mcball[of a ε]] mdiameter-cball-leq[of a ε]*
by *auto*

lemma *mdiameter-is-sup*:
assumes $x \in A \cap M \ y \in A \cap M$
shows $d \ x \ y \leq mdiameter \ A$


```

using assms by(auto simp: mdiameter-def intro!: Sup-upper)

lemma mdiameter-is-sup':
  assumes  $x \in A \cap M$   $y \in A \cap M$   $\text{mdiameter } A \leq \text{ennreal } r$   $r \geq 0$ 
  shows  $d \ x \ y \leq r$ 
  using order.trans[OF mdiameter-is-sup[OF assms(1,2)] assms(3)] assms(4) by
  simp

lemma mdiameter-le:
  assumes  $\bigwedge x \ y. x \in A \implies y \in A \implies d \ x \ y \leq r$ 
  shows  $\text{mdiameter } A \leq r$ 
  using assms by(auto simp: mdiameter-def Sup-le-iff ennreal-leI)

lemma mdiameter-eq-closure:  $\text{mdiameter } (\text{mtopology closure-of } A) = \text{mdiameter } A$ 
proof(rule antisym)
  show  $\text{mdiameter } A \leq \text{mdiameter } (\text{mtopology closure-of } A)$ 
  by(fastforce intro!: Sup-subset-mono simp: mdiameter-def metric-closure-of)
next
  have  $\{ \text{ennreal } (d \ x \ y) \mid x \ y. x \in A \cap M \wedge y \in A \cap M \} = \text{ennreal } \{ d \ x \ y \mid x \ y. x \in A \cap M \wedge y \in A \cap M \}$ 
  by auto
  also have  $\text{mdiameter } (\text{mtopology closure-of } A) \leq \bigsqcup \dots$ 
  unfolding le-Sup-iff-less
proof safe
  fix  $r$ 
  assume  $r < \text{mdiameter } (\text{mtopology closure-of } A)$ 
  then obtain  $x \ y$  where  $xy: x \in \text{mtopology closure-of } A$   $x \in M$   $y \in \text{mtopology closure-of } A$   $y \in M$   $r < \text{ennreal } (d \ x \ y)$ 
  by(auto simp: mdiameter-def less-Sup-iff)
  hence  $r < \top$ 
  using dual-order.strict-trans ennreal-less-top by blast
  define  $e$  where  $e \equiv (d \ x \ y - \text{enn2real } r) / 2$ 
  have  $e > 0$ 
  using  $xy(5)$   $\langle r < \top \rangle$  by(simp add: e-def)
  then obtain  $x' \ y'$  where  $xy': x' \in \text{mball } x \ e$   $x' \in A$   $y' \in \text{mball } y \ e$   $y' \in A$ 
  using  $xy$  by(fastforce simp: metric-closure-of)
  show  $\exists i \in \{ d \ x \ y \mid x \ y. x \in A \cap M \wedge y \in A \cap M \}. r \leq \text{ennreal } i$ 
  proof(safe intro!: bexI[where  $x = d \ x' \ y'$ ])
  have  $d \ x \ y \leq d \ x \ x' + d \ x' \ y' + d \ y \ y'$ 
  using triangle[OF  $xy(2) - xy(4)$ , of  $x'$ ]  $xy'$  triangle[of  $x' \ y' \ y$ ]
  by(fastforce simp add: commute)
  also have  $\dots < d \ x \ y - \text{enn2real } r + d \ x' \ y'$ 
  using  $xy'(1)$   $xy'(3)$  by(simp add: e-def)
  finally have  $\text{enn2real } r < d \ x' \ y'$  by simp
  thus  $r \leq \text{ennreal } (d \ x' \ y')$ 
  by (simp add:  $\langle r < \top \rangle$ )
qed(use  $xy'(1)$   $xy'(3)$   $xy'(2,4)$  in auto)
qed

```

```

finally show mdiameter (mtopology closure-of A) ≤ mdiameter A
  by(simp add: mdiameter-def)
qed

lemma mbounded-finite-mdiameter: mbounded A ↔ A ⊆ M ∧ mdiameter A <
∞
proof safe
  assume mbounded A
  then obtain x B where A ⊆ mball x B
    by(auto simp: mbounded-def)
  then have mdiameter A ≤ mdiameter (mball x B)
    by(rule mdiameter-subset)
  also have ... ≤ ennreal (2 * B)
    by(rule mdiameter-cball-leq)
  also have ... < ∞
    by auto
  finally show mdiameter A < ∞ .
next
  assume h:mdiameter A < ∞ A ⊆ M
  consider A = {} | A ≠ {} by auto
  then show mbounded A
  proof cases
    case h2:2
      then have 1:{d x y | x y. x ∈ A ∧ y ∈ A} ≠ {} by auto
      have eq:{ennreal (d x y) | x y. x ∈ A ∧ y ∈ A} = ennreal ‘ {d x y | x y. x ∈
A ∧ y ∈ A}
        by auto
      hence 2:mdiameter A = ⌊ (ennreal ‘ {d x y | x y. x ∈ A ∧ y ∈ A})
        using h by(auto simp add: mdiameter-def2)
      obtain x y where hxy:
        x ∈ A y ∈ A mdiameter A < ennreal (d x y + 1)
        using SUP-approx-ennreal[OF - 1 2,of 1] h by(fastforce simp: diameter-def)
      show ?thesis
        unfolding mbounded-alt
      proof(safe intro!: exI[where x=d x y + 1])
        fix w z
        assume w ∈ A z ∈ A
        with SUP-lessD[OF hxy(3)][simplified 2]]
        have ennreal (d w z) < ennreal (d x y + 1)
          by blast
        thus d w z ≤ d x y + 1
          by (metis canonically-ordered-monoid-add-class.lessE ennreal-le-iff2 en-
nreal-neg le-iff-add not-less-zero)
        qed (use h in auto)
      qed(auto simp: mbounded-def)
    qed(auto simp: mbounded-def)

```

Distance between a point and a set.

definition d-set :: 'a set ⇒ 'a ⇒ real **where**

$d\text{-set } A \equiv (\lambda x. \text{if } A \neq \{\} \wedge A \subseteq M \wedge x \in M \text{ then } \text{Inf } \{d\ x\ y \mid y. y \in A\} \text{ else } 0)$

lemma $d\text{-set-nonneg}[simp]$:

$d\text{-set } A\ x \geq 0$

proof –

have $\{d\ x\ y \mid y. y \in A\} = d\ x\ \text{' } A$ **by** *auto*

thus *?thesis*

by(*auto simp: d-set-def intro!: cINF-greatest[of - - d x]*)

qed

lemma $d\text{-set-bdd-below}[simp]$:

$\text{bdd-below } \{d\ x\ y \mid y. y \in A\}$

by(*auto simp: bdd-below-def intro!: exI[where x=0]*)

lemma $d\text{-set-singleton}[simp]$:

$x \in M \implies y \in M \implies d\text{-set } \{y\}\ x = d\ x\ y$

by(*auto simp: d-set-def*)

lemma $d\text{-set-empty}[simp]$:

$d\text{-set } \{\}\ x = 0$

by(*simp add: d-set-def*)

lemma $d\text{-set-notin}$:

$x \notin M \implies d\text{-set } A\ x = 0$

by(*auto simp: d-set-def*)

lemma $d\text{-set-inA}$:

assumes $x \in A$

shows $d\text{-set } A\ x = 0$

proof –

{

assume $x \in M\ A \subseteq M$

then have $0 \in \{d\ x\ y \mid y. y \in A\}$

using *assms* **by**(*auto intro: exI[where x=x]*)

moreover have $A \neq \{\}$

using *assms* **by** *auto*

ultimately have $\text{Inf } \{d\ x\ y \mid y. y \in A\} = 0$

using *cInf-lower[OF <0 ∈ {d x y | y. y ∈ A}>] d-set-nonneg[of A x] <A ⊆ M>*

$\langle x \in M \rangle$

by(*auto simp: d-set-def*)

}

thus *?thesis*

using *assms* **by**(*auto simp: d-set-def*)

qed

lemma $d\text{-set-nzeroD}$:

assumes $d\text{-set } A\ x \neq 0$

shows $A \subseteq M\ x \notin A\ A \neq \{\}$

by(*rule ccontr, use assms d-set-inA d-set-def in auto*)

lemma *d-set-antimono*:
assumes $A \subseteq B$ $A \neq \{\}$ $B \subseteq M$
shows $d\text{-set } B \ x \leq d\text{-set } A \ x$
proof (*cases* $B = \{\}$)
 case *h:False*
 with *assms* **have** $\{d \ x \ y \mid y. y \in B\} \neq \{\}$ $\{d \ x \ y \mid y. y \in A\} \subseteq \{d \ x \ y \mid y. y \in B\}$
 $A \subseteq M$
 by *auto*
 with *assms*(3) **show** ?thesis
 by(*simp add: d-set-def cInf-superset-mono assms*(2))
qed(*use assms in simp*)

lemma *d-set-bounded*:
assumes $\bigwedge y. y \in A \implies d \ x \ y < K$ $K > 0$
shows $d\text{-set } A \ x < K$
proof –
 consider $A = \{\} \mid \neg A \subseteq M \mid x \notin M \mid A \neq \{\} \mid A \subseteq M \mid x \in M$
 by *blast*
 then show ?thesis
 proof *cases*
 case 4
 then have $2: \{d \ x \ y \mid y. y \in A\} \neq \{\}$ **by** *auto*
 show ?thesis
 using *assms* **by** (*auto simp add: d-set-def cInf-lessD[OF 2] cInf-less-iff[OF 2]*)
qed(*auto simp: d-set-def assms*)
qed

lemma *d-set-tr*:
assumes $x \in M$ $y \in M$
shows $d\text{-set } A \ x \leq d \ x \ y + d\text{-set } A \ y$
proof –
 consider $A = \{\} \mid \neg A \subseteq M \mid A \neq \{\} \mid A \subseteq M$
 by *blast*
 then show ?thesis
 proof *cases*
 case 3
 have $d\text{-set } A \ x \leq \text{Inf } \{d \ x \ y + d \ y \ a \mid a. a \in A\}$
 proof –
 have $\bigcap \{d \ x \ y \mid y. y \in A\} \leq \bigcap \{d \ x \ y + d \ y \ a \mid a. a \in A\}$
 by(*rule cInf-mono*) (*use 3 assms triangle in fastforce*)+
 thus ?thesis
 by(*simp add: d-set-def assms 3*)
qed
 also have $\dots = (\bigcap a \in A. d \ x \ y + d \ y \ a)$
 by (*simp add: Setcompr-eq-image*)
 also have $\dots = d \ x \ y + \bigcap (d \ y \ ' A)$
 using 3 **by**(*auto intro!: Inf-add-eq bdd-belowI[where m=0]*)

also have $\dots = d\ x\ y + d\text{-set}\ A\ y$
using *assms* \mathfrak{I} **by** (*auto simp: d-set-def Setcompr-eq-image*)
finally show *?thesis* .
qed (*auto simp: d-set-def*)
qed

lemma *d-set-abs-le*:
assumes $x \in M\ y \in M$
shows $|d\text{-set}\ A\ x - d\text{-set}\ A\ y| \leq d\ x\ y$
using *d-set-tr[OF assms, of A] d-set-tr[OF assms(2,1), of A, simplified commute[of y x]]*
by *auto*

lemma *d-set-inA-le*:
assumes $y \in A$
shows $d\text{-set}\ A\ x \leq d\ x\ y$
proof –
consider $A \neq \{\}$ $A \subseteq M\ x \in M \mid \neg A \subseteq M \mid x \notin M$
using *assms* **by** *blast*
then show *?thesis*
proof *cases*
case 1
with *assms* **have** $y \in M$ **by** *auto*
from *d-set-tr[OF 1(3) this, of A]* **show** *?thesis*
by (*simp add: d-set-inA[OF assms]*)
qed (*auto simp: d-set-def*)
qed

lemma *d-set-ball-empty*:
assumes $A \neq \{\}$ $A \subseteq M\ e > 0\ x \in M\ mball\ x\ e \cap A = \{\}$
shows $d\text{-set}\ A\ x \geq e$
using *assms* **by** (*fastforce simp: d-set-def assms(1) le-cInf-iff*)

lemma *d-set-closed-pos*:
assumes *closedin mtopology A A* $A \neq \{\}$ $x \in M\ x \notin A$
shows $d\text{-set}\ A\ x > 0$
proof –
have $a: A \subseteq M$ *openin mtopology (M - A)*
using *closedin-subset[OF assms(1)] assms(1)* **by** *auto*
with *assms(3,4)* **obtain** e **where** $e > 0\ mball\ x\ e \subseteq M - A$
using *openin-mtopology* **by** *auto*
thus *?thesis*
by (*auto intro!: order.strict-trans2[OF e(1) d-set-ball-empty[OF assms(2) a(1) e(1) assms(3)]]*)
qed

lemma *gdelta-in-closed*:
assumes *closedin mtopology M*
shows *gdelta-in mtopology M*

by(*auto intro!*: *closed-imp-gdelta-in metrizable-space-mtopology*)

Oscillation

definition *osc-on* :: [*'b set, 'b topology, 'b \Rightarrow 'a, 'b*] \Rightarrow *ennreal* **where**
osc-on *A X f* \equiv ($\lambda y. \sqcap \{ \text{mdiameter } (f \text{ ' } (A \cap U)) \mid U. y \in U \wedge \text{openin } X U \}$)

abbreviation *osc X* \equiv *osc-on* (*topspace X*) *X*

lemma *osc-def*: *osc X f* = ($\lambda y. \sqcap \{ \text{mdiameter } (f \text{ ' } U) \mid U. y \in U \wedge \text{openin } X U \}$)

by(*standard, auto simp: osc-on-def*) (*metis (no-types, opaque-lifting) inf.absorb2 openin-subset*)

lemma *osc-on-less-iff*:

osc-on A X f x < t \longleftrightarrow ($\exists v. x \in v \wedge \text{openin } X v \wedge \text{mdiameter } (f \text{ ' } (A \cap v)) < t$)

by(*auto simp add: osc-on-def Inf-less-iff*)

lemma *osc-less-iff*:

osc X f x < t \longleftrightarrow ($\exists v. x \in v \wedge \text{openin } X v \wedge \text{mdiameter } (f \text{ ' } v) < t$)

by(*auto simp add: osc-def Inf-less-iff*)

end

definition *mdist-set* :: *'a metric \Rightarrow 'a set \Rightarrow 'a \Rightarrow real* **where**

mdist-set m \equiv *Metric-space.d-set* (*mspace m*) (*mdist m*)

lemma(*in Metric-space*) *mdist-set-Self*: *mdist-set Self* = *d-set*

by(*simp add: mdist-set-def*)

lemma *mdist-set-nonneg*[*simp*]: *mdist-set m A x* ≥ 0

by(*auto simp: mdist-set-def Metric-space.d-set-nonneg*)

lemma *mdist-set-singleton*[*simp*]:

$x \in \text{mspace } m \Longrightarrow y \in \text{mspace } m \Longrightarrow \text{mdist-set } m \{y\} x = \text{mdist } m x y$

by(*auto simp: mdist-set-def Metric-space.d-set-singleton*)

lemma *mdist-set-empty*[*simp*]: *mdist-set m {} x* = 0

by(*auto simp: mdist-set-def Metric-space.d-set-empty*)

lemma *mdist-set-inA*:

assumes $x \in A$

shows *mdist-set m A x* = 0

by(*auto simp: assms mdist-set-def Metric-space.d-set-inA*)

lemma *mdist-set-nzeroD*:

assumes *mdist-set m A x* $\neq 0$

shows $A \subseteq \text{mspace } m \ x \notin A \ A \neq \{\}$

using *assms Metric-space.d-set-nzeroD*[*of mspace m mdist m*]

by(*auto simp: mdist-set-def*)

lemma *mdist-set-antimono*:
assumes $A \subseteq B$ $A \neq \{\}$ $B \subseteq \text{mspace } m$
shows $\text{mdist-set } m B x \leq \text{mdist-set } m A x$
by(*auto simp: assms mdist-set-def Metric-space.d-set-antimono*)

lemma *mdist-set-bounded*:
assumes $\bigwedge y. y \in A \implies \text{mdist } m x y < K$ $K > 0$
shows $\text{mdist-set } m A x < K$
by(*auto simp: assms mdist-set-def Metric-space.d-set-bounded*)

lemma *mdist-set-tr*:
assumes $x \in \text{mspace } m$ $y \in \text{mspace } m$
shows $\text{mdist-set } m A x \leq \text{mdist } m x y + \text{mdist-set } m A y$
by(*auto simp: assms mdist-set-def Metric-space.d-set-tr*)

lemma *mdist-set-abs-le*:
assumes $x \in \text{mspace } m$ $y \in \text{mspace } m$
shows $|\text{mdist-set } m A x - \text{mdist-set } m A y| \leq \text{mdist } m x y$
by(*auto simp: assms mdist-set-def Metric-space.d-set-abs-le*)

lemma *mdist-set-inA-le*:
assumes $y \in A$
shows $\text{mdist-set } m A x \leq \text{mdist } m x y$
by(*auto simp: assms mdist-set-def Metric-space.d-set-inA-le*)

lemma *mdist-set-ball-empty*:
assumes $A \neq \{\}$ $A \subseteq \text{mspace } m$ $e > 0$ $x \in \text{mspace } m$ $\text{mball-of } m x e \cap A = \{\}$
shows $\text{mdist-set } m A x \geq e$
by (*metis Metric-space.d-set-ball-empty Metric-space-mspace-mdist assms mball-of-def mdist-set-def*)

lemma *mdist-set-closed-pos*:
assumes $\text{closedin } (\text{mtopology-of } m) A$ $A \neq \{\}$ $x \in \text{mspace } m$ $x \notin A$
shows $\text{mdist-set } m A x > 0$
by (*metis Metric-space.d-set-closed-pos Metric-space-mspace-mdist assms mdist-set-def mtopology-of-def*)

lemma *mdist-set-uniformly-continuous*: *uniformly-continuous-map* m *euclidean-metric* (*mdist-set* $m A$)
unfolding *uniformly-continuous-map-def*
proof *safe*
fix $e :: \text{real}$
assume $0 < e$
then show $\exists \delta > 0. \forall x \in \text{mspace } m. \forall y \in \text{mspace } m. \text{mdist } m y x < \delta \implies \text{mdist euclidean-metric } (\text{mdist-set } m A y) (\text{mdist-set } m A x) < e$
using *order.strict-trans1[OF mdist-set-abs-le]* **by**(*auto intro!: exI[where x=e] simp: dist-real-def*)
qed *simp*

lemma *uniformly-continuous-map-add*:
fixes $f :: 'a \Rightarrow 'b::\text{real-normed-vector}$
assumes *uniformly-continuous-map* m *euclidean-metric* f *uniformly-continuous-map*
m euclidean-metric g
shows *uniformly-continuous-map* m *euclidean-metric* $(\lambda x. f x + g x)$
unfolding *uniformly-continuous-map-def*
proof *safe*
fix $e :: \text{real}$
assume $e > 0$
from *half-gt-zero*[*OF this*] *assms* **obtain** $d1 d2$ **where** $d: d1 > 0 d2 > 0$
 $\bigwedge x y. x \in \text{mspace } m \implies y \in \text{mspace } m \implies \text{mdist } m x y < d1 \implies \text{dist } (f x) (f y) < e / 2$ $\bigwedge x y. x \in \text{mspace } m \implies y \in \text{mspace } m \implies \text{mdist } m x y < d2 \implies \text{dist } (g x) (g y) < e / 2$
by(*simp only: uniformly-continuous-map-def mdist-euclidean-metric*) *metis*
show $\exists \delta > 0. \forall y \in \text{mspace } m. \forall x \in \text{mspace } m. \text{mdist } m x y < \delta \longrightarrow \text{mdist euclidean-metric } (f x + g x) (f y + g y) < e$
using d **by**(*fastforce intro!: exI[where x=min d1 d2] order.strict-trans1[OF dist-triangle-add]*)
qed *simp*

lemma *uniformly-continuous-map-real-divide*:
fixes $f :: 'a \Rightarrow \text{real}$
assumes *uniformly-continuous-map* m *euclidean-metric* f *uniformly-continuous-map*
m euclidean-metric g
and $\bigwedge x. x \in \text{mspace } m \implies g x \neq 0$ $\bigwedge x. x \in \text{mspace } m \implies |g x| \geq a a > 0$
 $\bigwedge x. x \in \text{mspace } m \implies |g x| < Kg$
and $\bigwedge x. x \in \text{mspace } m \implies |f x| < Kf$
shows *uniformly-continuous-map* m *euclidean-metric* $(\lambda x. f x / g x)$
unfolding *uniformly-continuous-map-def*
proof *safe*
fix $e :: \text{real}$
assume $e[\text{arith}]: e > 0$
consider $\text{mspace } m \neq \{\}$ | $\text{mspace } m = \{\}$ **by** *auto*
then show $\exists \delta > 0. \forall x \in \text{mspace } m. \forall y \in \text{mspace } m. \text{mdist } m y x < \delta \longrightarrow \text{mdist euclidean-metric } (f y / g y) (f x / g x) < e$
proof *cases*
case $m:1$
then have $Kfg\text{-pos}[\text{arith}]: Kg > 0 Kf \geq 0$
using $\text{assms}(4-7)$ **by** *auto fastforce+*
define e' **where** $e' \equiv a^2 / (Kf + Kg) * e$
have $e':e' > 0$
using $\text{assms}(5)$ **by**(*auto simp: e'-def*)
with assms **obtain** $d1 d2$ **where** $d: d1 > 0 d2 > 0$
 $\bigwedge x y. x \in \text{mspace } m \implies y \in \text{mspace } m \implies \text{mdist } m x y < d1 \implies |f x - f y| < e'$ $\bigwedge x y. x \in \text{mspace } m \implies y \in \text{mspace } m \implies \text{mdist } m x y < d2 \implies |g x - g y| < e'$
by(*simp only: uniformly-continuous-map-def mdist-euclidean-metric dist-real-def*) *metis*
show *?thesis*


```

unfolding mdist-euclidean-metric dist-real-def
proof(safe intro!: exI[where  $x = \min d1 d2$ ])
  fix  $x y$ 
  assume  $x : x \in \text{mspace } m$  and  $y : y \in \text{mspace } m$  and  $\text{mdist } m x y < \min d1 d2$ 
  then have  $\text{dist}[\text{arith}] : \text{mdist } m x y < d1$   $\text{mdist } m x y < d2$  by auto
  note  $[\text{arith}] = \text{assms}(3,4,6,7)[OF x]$   $\text{assms}(3,4,6,7)[OF y]$ 
  have  $|f x / g x - f y / g y| = |(f x * g y - f y * g x) / (g x * g y)|$ 
    by(simp add: diff-frac-eq)
  also have  $\dots = |(f x * g y - f x * g x + (f x * g x - f y * g x)) / (g x * g y)|$ 
    by simp
  also have  $\dots = |(f x - f y) * g x - f x * (g x - g y)| / (|g x| * |g y|)$ 
    by(simp add: left-diff-distrib right-diff-distrib abs-mult)
  also have  $\dots \leq (|f x - f y| * |g x| + |f x| * |g x - g y|) / (|g x| * |g y|)$ 
    by(rule divide-right-mono) (use abs-triangle-ineq4 abs-mult inmetis,auto)
  also have  $\dots < (e' * Kg + Kf * e') / (|g x| * |g y|)$ 
    by(rule divide-strict-right-mono[OF add-less-le-mono]) (auto intro! : mult-mono'
mult-strict-mono, use d(3,4)[OF x y] in auto)
  also have  $\dots \leq (e' * Kg + Kf * e') / a^{\wedge}2$ 
    by(auto intro! : divide-left-mono simp: power2-eq-square) (insert assms(5)
e', auto simp: <a ≤ |g x|> mult-mono')
  also have  $\dots = (Kf + Kg) / a^{\wedge}2 * e'$ 
    by (simp add: distrib-left mult.commute)
  also have  $\dots = e$ 
    using assms(5) by(auto simp: e'-def)
  finally show  $|f x / g x - f y / g y| < e$  .
  qed(use d in auto)
qed (auto intro! : exI[where  $x=1$ ])
qed simp

```

lemma

```

assumes  $e > 0$ 
shows uniformly-continuous-map-from-capped-metric:uniformly-continuous-map
(capped-metric e m1)  $m2 f \longleftrightarrow$  uniformly-continuous-map m1 m2 f (is ?g1)
and uniformly-continuous-map-to-capped-metric:uniformly-continuous-map m1
(capped-metric e m2)  $f \longleftrightarrow$  uniformly-continuous-map m1 m2 f (is ?g2)
proof -
have [simp]:( $\lambda n. \min e (X n) \longrightarrow 0 \longleftrightarrow X \longrightarrow 0$  for  $X$ )
proof
assume  $h : (\lambda n. \min e (X n) \longrightarrow 0$ 
show  $X \longrightarrow 0$ 
unfolding LIMSEQ-def dist-real-def
proof safe
fix  $r :: \text{real}$ 
assume  $0 < r$ 
then have  $\min (e / 2) r > 0$ 
using assms by auto
from LIMSEQ-D[OF h this] obtain  $N$  where  $N : \bigwedge n. n \geq N \implies |\min e (X n)| < \min (e / 2) r$ 
by auto

```

```

have  $\min e (X n) = X n$  if  $h:n \geq N$  for  $n$ 
proof(rule ccontr)
  assume  $\min e (X n) \neq X n$ 
  then have  $\min e (X n) = e$ 
    by auto
  with  $N[OF h]$  show False
    by auto
qed
with  $N$  show  $\exists no. \forall n \geq no. |X n - 0| < r$ 
  by(auto intro!: exI[where  $x=N$ ])
qed
qed(auto intro!: tendsto-eq-intros(4)[of  $\lambda x. e$  e sequentially  $X -$ ] simp: assms)
show ?g1 ?g2
  using assms by(auto simp: uniformly-continuous-map-sequentially capped-metric-mdist)
qed

```

lemma *Urysohn-lemma-uniform:*

```

assumes closedin (mtopology-of  $m$ )  $T$  closedin (mtopology-of  $m$ )  $U$   $T \cap U = \{\}$ 
 $\bigwedge x y. x \in T \implies y \in U \implies \text{mdist } m x y \geq e$   $e > 0$ 

```

```

obtains  $f :: 'a \Rightarrow \text{real}$ 

```

```

where uniformly-continuous-map  $m$  euclidean-metric  $f$ 

```

```

 $\bigwedge x. f x \geq 0$   $\bigwedge x. f x \leq 1$   $\bigwedge x. x \in T \implies f x = 1$   $\bigwedge x. x \in U \implies f x = 0$ 

```

```

proof -

```

```

consider  $T = \{\}$  |  $U = \{\}$  |  $T \neq \{\}$   $U \neq \{\}$  by auto

```

```

then show ?thesis

```

```

proof cases

```

```

  case 1

```

```

    define  $f$  where  $f \equiv (\lambda x::'a. 0::\text{real})$ 

```

```

    with 1 have uniformly-continuous-map  $m$  euclidean-metric  $f$   $\bigwedge x. f x \geq 0$   $\bigwedge x.$ 
 $f x \leq 1$   $\bigwedge x. x \in T \implies f x = 1$   $\bigwedge x. x \in U \implies f x = 0$ 

```

```

      by auto

```

```

    then show ?thesis

```

```

      using that by auto

```

```

  next

```

```

    case 2

```

```

    define  $f$  where  $f \equiv (\lambda x::'a. 1::\text{real})$ 

```

```

    with 2 have uniformly-continuous-map  $m$  euclidean-metric  $f$   $\bigwedge x. f x \geq 0$   $\bigwedge x.$ 
 $f x \leq 1$   $\bigwedge x. x \in T \implies f x = 1$   $\bigwedge x. x \in U \implies f x = 0$ 

```

```

      by auto

```

```

    then show ?thesis

```

```

      using that by auto

```

```

  next

```

```

    case  $TU:3$ 

```

```

    then have  $STU:\text{mspace } m \neq \{\}$   $T \subseteq \text{mspace } m$   $U \subseteq \text{mspace } m$ 

```

```

      using assms(1,2) closedin-topospace-empty closedin-subset by fastforce+

```

```

    interpret  $m$ : Metric-space  $\text{mspace } m$  mdist (capped-metric  $e$   $m$ )

```

```

      by (metis Metric-space-mspace-mdist capped-metric-mspace)

```

```

    have  $e:x \in T \implies y \in U \implies \text{mdist} (\text{capped-metric } e m) x y \geq e$  for  $x y$ 

```

```

      by (simp add: assms(4) capped-metric-mdist)

```

```

define f where  $f \equiv (\lambda x. \text{mdist-set } (\text{capped-metric } e \ m) \ U \ x \ / \ (\text{mdist-set } (\text{capped-metric } e \ m) \ U \ x \ + \ \text{mdist-set } (\text{capped-metric } e \ m) \ T \ x))$ 
have uniformly-continuous-map m euclidean-metric f
unfolding uniformly-continuous-map-from-capped-metric[OF assms(5),of
m,symmetric] f-def
proof(rule uniformly-continuous-map-real-divide[where  $Kf=e + 1$  and  $Kg=2 * e + 1$  and  $a=e / 2$ ])
show uniformly-continuous-map (capped-metric e m) euclidean-metric (mdist-set
(capped-metric e m) U)
uniformly-continuous-map (capped-metric e m) euclidean-metric ( $\lambda x.$ 
mdist-set (capped-metric e m) U x + mdist-set (capped-metric e m) T x)
by(auto intro!: mdist-set-uniformly-continuous uniformly-continuous-map-add)
next
fix x
assume  $x : x \in \text{mspace } (\text{capped-metric } e \ m)$ 
then consider  $x \notin T \mid x \notin U$ 
using TU assms by auto
thus mdist-set (capped-metric e m) U x + mdist-set (capped-metric e m) T
 $x \neq 0$ 
proof cases
case 1
with TU x assms have mdist-set (capped-metric e m) T  $x > 0$ 
by(auto intro!: mdist-set-closed-pos simp: mtopology-capped-metric)
thus ?thesis
by (metis add-less-same-cancel2 linorder-not-less mdist-set-nonneg)
next
case 2
with TU x assms have mdist-set (capped-metric e m) U  $x > 0$ 
by(auto intro!: mdist-set-closed-pos simp: mtopology-capped-metric)
thus ?thesis
by (metis add-less-same-cancel1 linorder-not-less mdist-set-nonneg)
qed
next
fix x
assume  $x : x \in \text{mspace } (\text{capped-metric } e \ m)$ 
consider  $x \in (\bigcup a \in U. \text{m.ball } a \ (e / 2)) \mid x \notin (\bigcup a \in U. \text{m.ball } a \ (e / 2))$ 
by auto
then show  $e / 2 \leq \text{mdist-set } (\text{capped-metric } e \ m) \ U \ x \ + \ \text{mdist-set } (\text{capped-metric } e \ m) \ T \ x$ 
proof cases
case 1
have  $\text{m.ball } x \ (e / 2) \cap T = \{\}$ 
proof(rule ccontr)
assume  $\text{m.ball } x \ (e / 2) \cap T \neq \{\}$ 
then obtain y where  $y : y \in \text{m.ball } x \ (e / 2) \ y \in T$ 
by blast
then obtain u where  $u : u \in U \ x \in \text{m.ball } u \ (e / 2)$ 
using 1 by auto
have  $\text{mdist } (\text{capped-metric } e \ m) \ y \ u \leq \text{mdist } (\text{capped-metric } e \ m) \ y \ x \ +$ 

```

```

mdist (capped-metric e m) x u
  using STU u y x by(auto intro!: m.triangle)
  also have ... < e / 2 + e / 2
  using y u by(auto simp: m commute)
  also have ... = e using assms(5) by linarith
  finally show False
  using e[OF y(2) u(1)] by simp
qed
from m.d-set-ball-empty[OF - - - this] TU STU x assms(1,5)
have e / 2 ≤ m.d-set T x
  using STU(2) x by auto
also have ... ≤ |mdist-set (capped-metric e m) U x + mdist-set (capped-metric
e m) T x|
  by (simp add: mdist-set-def)
  finally show ?thesis .
next
case 2
then have m.mball x (e / 2) ∩ U = {}
  using x by (auto simp: m commute)
hence e / 2 ≤ m.d-set U x
  by (metis STU(3) TU(2) assms(5) capped-metric-mspace half-gt-zero
m.d-set-ball-empty x)
also have ... ≤ |mdist-set (capped-metric e m) U x + mdist-set (capped-metric
e m) T x|
  by (simp add: mdist-set-def)
  finally show ?thesis .
qed
next
fix x
assume x ∈ mspace (capped-metric e m)
have |mdist-set (capped-metric e m) U x + mdist-set (capped-metric e m) T
x| = mdist-set (capped-metric e m) U x + mdist-set (capped-metric e m) T x
  using mdist-set-nonneg by auto
also have ... < (e + 1 / 2) + (e + 1 / 2)
  apply(intro add-strict-mono mdist-set-bounded)
  using assms(5) add-strict-increasing[of 1 / 2, OF - mdist-capped[OF
assms(5), of m x]] by (auto simp add: add commute)
also have ... = 2 * e + 1
  by auto
finally show |mdist-set (capped-metric e m) U x + mdist-set (capped-metric
e m) T x| < 2 * e + 1 .
show |mdist-set (capped-metric e m) U x| < e + 1
  using assms(5) add-strict-increasing[of 1, OF - mdist-capped[OF assms(5), of
m x]] by (auto simp add: add commute intro!: mdist-set-bounded)
qed(use assms in auto)
moreover have ∧x. f x ∈ {0..1}
proof -
fix x
have f x ≤ mdist-set (capped-metric e m) U x / mdist-set (capped-metric e

```

```

m) U x
  proof –
    consider mdist-set (capped-metric e m) U x = 0 | mdist-set (capped-metric
e m) U x > 0
    using antisym-conv1 by fastforce
    thus ?thesis
    proof cases
      case 2
        show ?thesis
          unfolding f-def by(rule divide-left-mono[OF - - mult-pos-pos]) (auto
simp: 2 add-strict-increasing)
          qed (simp add: f-def)
        qed
        also have ... ≤ 1 by simp
        finally show f x ∈ {0..1}
          by (auto simp: f-def)
        qed
        moreover have f x = 1 if x:x ∈ T for x
        proof –
          { assume h:mdist-set (capped-metric e m) U x = 0
            then have x ∉ U using assms STU x by blast
            hence False
            by (metis STU(2) TU(2) assms(2) capped-metric-mspace h mdist-set-closed-pos
mtopology-capped-metric order-less-irrefl subset-eq x)
          }
          thus ?thesis
            by (metis add.right-neutral divide-self-if f-def mdist-set-nzeroD(2) x)
          qed
        moreover have ∧x. x ∈ U ⇒ f x = 0
          by (simp add: f-def m.d-set-inA mdist-set-def)
        ultimately show ?thesis
          using that by auto
        qed
      qed
    qed

```

Open maps

lemma *Metric-space-open-map-from-dist:*

```

  assumes f ∈ mspace m1 → mspace m2
  and ∧x ε. x ∈ mspace m1 ⇒ ε > 0 ⇒ ∃δ>0. ∀y∈mspace m1. mdist m2
(f x) (f y) < δ → mdist m1 x y < ε
  shows open-map (mtopology-of m1) (subtopology (mtopology-of m2) (f ‘ mspace
m1)) f

```

proof –

```

  interpret m1: Metric-space mspace m1 mdist m1 by simp
  interpret m2: Metric-space mspace m2 mdist m2 by simp
  interpret m2': Submetric mspace m2 mdist m2 f ‘ mspace m1
  using assms(1) by(auto simp: Submetric-def Submetric-axioms-def)
  have open-map m1.mtopology m2'.sub.mtopology f
  proof (safe intro!: open-map-with-base[OF m1.mtopology-base-in-balls])

```

```

fix a and e :: real
assume h:a ∈ mspace m1 0 < e
show openin m2'.sub.mtopology (f ' m1.mball a e)
  unfolding m2'.sub.openin-mtopology
proof safe
  fix x
  assume x:x ∈ m1.mball a e
  then have xs:x ∈ mspace m1
    by auto
  have 0 < e - mdist m1 a x
    using x by auto
  from assms(2)[OF xs this] obtain d where d:d > 0 ∧ y. y ∈ mspace m1
  ⇒ mdist m2 (f x) (f y) < d ⇒ mdist m1 x y < e - mdist m1 a x
    by auto
  show ∃ r>0. m2'.sub.mball (f x) r ⊆ f ' m1.mball a e
  proof(safe intro!: exI[where x=d])
    fix z
    assume z:z ∈ m2'.sub.mball (f x) d
    then obtain y where y:z = f y y ∈ mspace m1
      by auto
    hence mdist m2 (f x) (f y) < d
      using z by simp
    hence mdist m1 x y < e - mdist m1 a x
      using d y by auto
    hence mdist m1 a y < e
      using h(1) x y m1.triangle[of a x y] by auto
    with h(1) show z ∈ f ' m1.mball a e
      by(auto simp: y)
    qed fact
  qed auto
qed
thus open-map (mtopology-of m1) (subtopology (mtopology-of m2) (f ' mspace
m1)) f
  by (simp add: mtopology-of-def m2'.mtopology-submetric)
qed

```

1.3.3 Separability in Metric Spaces

```

context Metric-space
begin

```

For a metric space M , M is separable iff M is second countable.

```

lemma generated-by-countable-balls:

```

```

  assumes countable U and mdense U
  shows mtopology = topology-generated-by {mball y (1 / real n) | y n. y ∈ U}
proof -
  have hu: U ⊆ M ∧ x ε. x ∈ M ⇒ ε > 0 ⇒ mball x ε ∩ U ≠ {}
    using assms by(auto simp: mdense-def)
  show ?thesis

```

```

unfolding base-is-subbase[OF mtopology-base-in-balls,simplified subbase-in-def]
proof(rule topology-generated-by-eq)
  fix K
  assume K ∈ {mball y (1 / real n) | y n. y ∈ U}
  then show openin (topology-generated-by {mball a ε | a ε. a ∈ M ∧ 0 < ε}) K
  by(auto simp: base-is-subbase[OF mtopology-base-in-balls,simplified subbase-in-def,symmetric])
next
  have h0: ∧x ε. x ∈ M ⇒ ε > 0 ⇒ ∃y ∈ U. ∃n. x ∈ mball y (1 / real n) ∧
mball y (1 / real n) ⊆ mball x ε
  proof -
    fix x ε
    assume h: x ∈ M (0 :: real) < ε
    then obtain N where hn: 1 / ε < real N
      using reals-Archimedean2 by blast
    have hn0: 0 < real N
      by(rule ccontr) (use hn h in fastforce)
    hence hn': 1 / real N < ε
      using h hn by (metis divide-less-eq mult.commute)
    have mball x (1 / (2 * real N)) ∩ U ≠ {}
      using mdense-def[of U] assms(2) h(1) hn0 by fastforce
    then obtain y where hy:
      y ∈ U y ∈ M y ∈ mball x (1 / (real (2 * N)))
      using hu by auto
    show ∃y ∈ U. ∃n. x ∈ mball y (1 / real n) ∧ mball y (1 / real n) ⊆ mball x ε
    proof(intro bexI[where x=y] exI[where x=2 * N] conjI)
      show x ∈ mball y (1 / real (2 * N))
        using hy(3) by (auto simp: commute)
    next
    show mball y (1 / real (2 * N)) ⊆ mball x ε
    proof
      fix y'
      assume hy': y' ∈ mball y (1 / real (2 * N))
      have d x y' < ε (is ?lhs < ?rhs)
      proof -
        have ?lhs ≤ d x y + d y y'
          using hy(2) hy' h(1) triangle by auto
        also have ... < 1 / real (2 * N) + 1 / real (2 * N)
          by(rule strict-ordered-ab-semigroup-add-class.add-strict-mono) (use
hy(3) hy(2) hy' h(1) hy' in auto)
        finally show ?thesis
          using hn' by auto
      qed
      thus y' ∈ mball x ε
        using hy' h(1) by auto
    qed
  qed fact
qed
fix K
assume hk: K ∈ {mball a ε | a ε. a ∈ M ∧ 0 < ε}

```

then obtain $x \in x$ **where** hxe :
 $x \in M \ 0 < \varepsilon x \ K = \text{mball } x \ \varepsilon x$ **by** *auto*
have $gh:K = (\bigcup \{\text{mball } y \ (1 / \text{real } n) \mid y \ n. \ y \in U \wedge \text{mball } y \ (1 / \text{real } n) \subseteq K\})$
proof
show $K \subseteq (\bigcup \{\text{mball } y \ (1 / \text{real } n) \mid y \ n. \ y \in U \wedge \text{mball } y \ (1 / \text{real } n) \subseteq K\})$
proof
fix k
assume $hkink:k \in K$
then have $hkinS:k \in M$
by (*simp add: hxe*)
obtain εk **where** $hek: \varepsilon k > 0 \ \text{mball } k \ \varepsilon k \subseteq K$
by (*metis Metric-space.openin-mtopology Metric-space-axioms hxe(3) hkink openin-mball*)
obtain $y \ n$ **where** $hyey$:
 $y \in U \ k \in \text{mball } y \ (1 / \text{real } n) \ \text{mball } y \ (1 / \text{real } n) \subseteq \text{mball } k \ \varepsilon k$
using *h0[OF hkinS hek(1)] by auto*
show $k \in \bigcup \{\text{mball } y \ (1 / \text{real } n) \mid y \ n. \ y \in U \wedge \text{mball } y \ (1 / \text{real } n) \subseteq K\}$
using *hek(2) hyey by blast*
qed
qed *auto*
show *openin (topology-generated-by {mball y (1 / real n) | y n. y ∈ U}) K*
unfolding *openin-topology-generated-by-iff*
apply (*rule generate-topology-on.UN[of {mball y (1 / real n) | y n. y ∈ U ∧ mball y (1 / real n) ⊆ K}, simplified gh[symmetric]]*)
apply (*rule generate-topology-on.Basis*) **by** *auto*
qed
qed

lemma *separable-space-imp-second-countable*:
assumes *separable-space mtopology*
shows *second-countable mtopology*
proof –
obtain U **where** hu :
 $\text{countable } U \ \text{mdense } U$
using *assms separable-space-def2 by blast*
show *?thesis*
proof (*rule countable-base-from-countable-subbase[where $\mathcal{O} = \{\text{mball } y \ (1 / \text{real } n) \mid y \ n. \ y \in U\}$]*)
have $\{\text{mball } y \ (1 / \text{real } n) \mid y \ n. \ y \in U\} = (\lambda(y,n). \ \text{mball } y \ (1 / \text{real } n)) \ ' (U \times \text{UNIV})$
by *auto*
also have *countable ...*
using hu **by** *auto*
finally show *countable {mball y (1 / real n) | y n. y ∈ U} .*
qed (*use subbase-in-def generated-by-countable-balls[of U] hu in auto*)
qed

corollary *separable-space-iff-second-countable:*
separable-space mtopology \longleftrightarrow *second-countable mtopology*
using *second-countable-imp-separable-space separable-space-imp-second-countable*
by *auto*

lemma *Lindelof-mdiameter:*

assumes *separable-space mtopology* $0 < e$
shows $\exists U. \text{countable } U \wedge \bigcup U = M \wedge (\forall u \in U. \text{mdiameter } u < \text{ennreal } e)$
proof –
have $(\bigwedge u. u \in \{\text{mball } x (e / 3) \mid x. x \in M\} \implies \text{openin mtopology } u)$
by *auto*
moreover have $\bigcup \{\text{mball } x (e / 3) \mid x. x \in M\} = M$
using *assms by auto*
ultimately have $\exists U'. \text{countable } U' \wedge U' \subseteq \{\text{mball } x (e / 3) \mid x. x \in M\} \wedge \bigcup U' = M$
using *second-countable-imp-Lindelof-space[OF assms(1)[simplified separable-space-iff-second-countable]]*
by(*auto simp: Lindelof-space-def*)
then obtain U' **where** $U': \text{countable } U' \wedge U' \subseteq \{\text{mball } x (e / 3) \mid x. x \in M\} \wedge \bigcup U' = M$
by *auto*
show *?thesis*
proof(*safe intro!: exI[where x=U']*)
fix u
assume $u \in U'$
then obtain x **where** $u : u = \text{mball } x (e / 3)$
using U' **by** *auto*
have $\text{mdiameter } u \leq \text{ennreal } (2 * (e / 3))$
by(*simp only: u mdiameter-ball-leq*)
also have $\dots < \text{ennreal } e$
by(*auto intro!: ennreal-lessI assms*)
finally show $\text{mdiameter } u < \text{ennreal } e$.
qed(*use U' in auto*)
qed

end

lemma *metrizable-space-separable-iff-second-countable:*

assumes *metrizable-space X*
shows *separable-space X* \longleftrightarrow *second-countable X*
proof –
obtain d **where** *Metric-space (topspace X) d Metric-space.mtopology (topspace X) d = X*
by (*metis assms(1) Metric-space.topspace-mtopology metrizable-space-def*)
thus *?thesis*
using *Metric-space.separable-space-iff-second-countable by fastforce*
qed

abbreviation $\text{mdense-of } m \ U \equiv \text{dense-in (mtopology-of } m) \ U$

lemma *mdense-of-def*: $mdense\text{-of } m \ U \longleftrightarrow (U \subseteq mspace \ m \wedge (\forall x \in mspace \ m. \forall \varepsilon > 0. mball\text{-of } m \ x \ \varepsilon \cap U \neq \{\}))$

using *Metric-space.mdense-def*[of *mspace m mdist m*] **by** (*simp add: mball-of-def mtopology-of-def*)

lemma *mdense-of-def2*: $mdense\text{-of } m \ U \longleftrightarrow (U \subseteq mspace \ m \wedge (\forall x \in mspace \ m. \forall \varepsilon > 0. \exists y \in U. mdist \ m \ x \ y < \varepsilon))$

using *Metric-space.mdense-def2*[of *mspace m mdist m*] **by** (*simp add: mtopology-of-def*)

lemma *mdense-of-def3*: $mdense\text{-of } m \ U \longleftrightarrow (U \subseteq mspace \ m \wedge (\forall x \in mspace \ m. \exists xn. range \ xn \subseteq U \wedge limitin \ (mtopology\text{-of } m) \ xn \ x \ sequentially))$

using *Metric-space.mdense-def3*[of *mspace m mdist m*] **by** (*simp add: mtopology-of-def*)

1.3.4 Compact Metric Spaces

context *Metric-space*

begin

lemma *mtotally-bounded-eq-compact-closedin*:

assumes *mcomplete closedin mtopology S*

shows *mtotally-bounded S* $\longleftrightarrow S \subseteq M \wedge compactin \ mtopology \ S$

by (*metis assms closure-of-eq mtotally-bounded-eq-compact-closure-of*)

lemma *mtotally-bounded-def2*: *mtotally-bounded S* $\longleftrightarrow (\forall \varepsilon > 0. \exists K. finite \ K \wedge K \subseteq M \wedge S \subseteq (\bigcup x \in K. mball \ x \ \varepsilon))$

unfolding *mtotally-bounded-def*

proof *safe*

fix *e :: real*

assume *h:e > 0* $\forall \varepsilon > 0. \exists K. finite \ K \wedge K \subseteq S \wedge S \subseteq (\bigcup x \in K. mball \ x \ \varepsilon)$

then show $\exists K. finite \ K \wedge K \subseteq M \wedge S \subseteq (\bigcup x \in K. mball \ x \ e)$

by (*metis Metric-space.mbounded-subset Metric-space.mtotally-bounded-imp-mbounded Metric-space-axioms mbounded-subset-mspace mtotally-bounded-def*)

next

fix *e :: real*

assume *e > 0* $\forall \varepsilon > 0. \exists K. finite \ K \wedge K \subseteq M \wedge S \subseteq (\bigcup x \in K. mball \ x \ \varepsilon)$

then obtain *K* **where** *K: finite K K* $\subseteq M \wedge S \subseteq (\bigcup x \in K. mball \ x \ (e / 2))$

by (*meson half-gt-zero*)

define *K'* **where** *K' ≡ {x ∈ K. mball x (e / 2) ∩ S ≠ {}}*

have *K'1: finite K' K' ⊆ M*

using *K* **by** (*auto simp: K'-def*)

have *K'2: S ⊆ (⋃ x ∈ K'. mball x (e / 2))*

proof

fix *x*

assume *x: x ∈ S*

then obtain *k* **where** *k: k ∈ K x ∈ mball k (e / 2)*

using *K* **by** *auto*

with *x* **have** *k ∈ K'*

```

    by(auto simp: K'-def)
  with k show x ∈ (⋃ x∈K'. mball x (e / 2))
    by auto
qed
have ∀ k∈K'. ∃ y. y ∈ mball k (e / 2) ∩ S
  by(auto simp: K'-def)
then obtain xk where xk: ∧k. k ∈ K' ⇒ xk k ∈ mball k (e / 2) ∧ k. k ∈ K'
⇒ xk k ∈ S
  by (metis IntD2 inf-commute)
hence ∧k. k ∈ K' ⇒ mball k (e / 2) ⊆ mball (xk k) e
  using triangle commute by fastforce
hence (⋃ x∈K'. mball x (e / 2)) ⊆ (⋃ x∈xk ' K'. mball x e)
  by blast
with K'2 have S ⊆ (⋃ x∈xk ' K'. mball x e)
  by blast
thus ∃ K. finite K ∧ K ⊆ S ∧ S ⊆ (⋃ x∈K. mball x e)
  by(intro exI[where x=xk ' K]) (use xk(2) K'1(1) in blast)
qed

```

lemma *compact-space-imp-separable:*

assumes *compact-space mtopology*

shows *separable-space mtopology*

proof –

have $\exists A. \text{finite } A \wedge A \subseteq M \wedge M \subseteq \bigcup ((\lambda a. \text{mball } a (1 / \text{real } (\text{Suc } n))) ' A)$
for n

using *assms* **by**(auto simp: *compact-space-eq-mcomplete-mtotally-bounded mtotally-bounded-def*)

then obtain A **where** $A: \bigwedge n. \text{finite } (A n) \wedge n. A n \subseteq M \wedge n. M \subseteq \bigcup ((\lambda a. \text{mball } a (1 / \text{real } (\text{Suc } n))) ' (A n))$

by *metis*

define K **where** $K \equiv \bigcup (\text{range } A)$

have $1: \text{countable } K$

using $A(1)$ **by**(auto intro!: *countable-UN[of - id,simplified]* simp: *K-def countable-finite*)

show *separable-space mtopology*

unfolding *mdense-def2 separable-space-def2*

proof(safe intro!: *exI[where x=K]* 1)

fix x **and** $\varepsilon :: \text{real}$

assume $h: x \in M \ 0 < \varepsilon$

then obtain n **where** $n: 1 / \text{real } (\text{Suc } n) \leq \varepsilon$

by (*meson nat-approx-posE order.strict-iff-not*)

then obtain y **where** $y: y \in A n \ x \in \text{mball } y (1 / \text{real } (\text{Suc } n))$

using $h(1)$ $A(3)[\text{of } n]$ **by** *auto*

thus $\exists y \in K. d x y < \varepsilon$

using n **by**(auto intro!: *beXI[where x=y]* simp: *commute K-def*)

qed(use *K-def A(2)* in *auto*)

qed

lemma *separable-space-cfunspace:*

```

assumes separable-space mtopology mcomplete
  and metrizable-space X compact-space X
  shows separable-space (mtopology-of (cfunspace X Self))
proof –
  consider topspace X = {} | topspace X ≠ {} M = {} | topspace X ≠ {} M ≠
  {} by auto
  thus ?thesis
  proof cases
    case 1
    then show ?thesis
      by(auto intro!: countable-space-separable-space)
  next
    case 2
    then have [simp]:mtopology = trivial-topology
      using topspace-mtopology by auto
    have 1:topspace (mtopology-of (cfunspace X Self)) = {}
      apply simp
      using 2(1) by(auto simp: mtopology-of-def)
    show ?thesis
      by(rule countable-space-separable-space, simp only: 1) simp
  next
    case XS-nem:3
    have cm: mcomplete-of (cfunspace X Self)
      by (simp add: assms(2) mcomplete-cfunspace)
    show ?thesis
    proof –
      obtain dx where dx: Metric-space (topspace X) dx Metric-space.mtopology
      (topspace X) dx = X
      by (metis Metric-space.topspace-mtopology assms(3) metrizable-space-def)
      interpret dx: Metric-space topspace X dx
      by fact
      have dx-c: dx.mcomplete
        using assms by(auto intro!: dx.compact-space-imp-mcomplete simp: dx(2))
      have ∃ B. finite B ∧ B ⊆ topspace X ∧ topspace X ⊆ (⋃ a∈B. dx.mball a (1
      / Suc m)) for m
        using dx.compactin-imp-mtotally-bounded[of topspace X] assms(4)
        by(fastforce simp: dx(2) compact-space-def dx.mtotally-bounded-def2)
      then obtain Xm where Xm: ∧ m. finite (Xm m) ∧ m. (Xm m) ⊆ topspace
      X ∧ m. topspace X ⊆ (⋃ a∈Xm m. dx.mball a (1 / Suc m))
        by metis
      hence Xm-eq: ∧ m. topspace X = (⋃ a∈Xm m. dx.mball a (1 / Suc m))
        by fastforce
      have Xm-nem: ∧ m. (Xm m) ≠ {}
        using XS-nem Xm-eq by blast
      define xmk where xmk ≡ (λ m. from-nat-into (Xm m))
      define km where km ≡ (λ m. card (Xm m))
      have km-pos: km m > 0 for m
        by(simp add: km-def card-gt-0-iff Xm Xm-nem)
      have xmk-bij: bij-betw (xmk m) {..for m

```

```

using bij-betw-from-nat-into-finite[OF Xm(1)] by (simp add: km-def xmk-def)
have xmk-into:  $xmk\ m\ i \in Xm\ m$  for  $m\ i$ 
  by (simp add: Xm-nem from-nat-into xmk-def)
have  $\exists U. countable\ U \wedge \bigcup U = M \wedge (\forall u \in U. mdiameter\ u < 1 / (Suc\ l))$ 
for  $l$ 
  by (rule Lindelof-mdiameter[OF assms(1)]) auto
  then obtain  $U$  where  $U: \bigwedge l. countable\ (U\ l) \wedge l. (\bigcup (U\ l)) = M \wedge l\ u. u$ 
 $\in U\ l \implies mdiameter\ u < 1 / Suc\ l$ 
  by metis
  have Ul-nem:  $U\ l \neq \{\}$  for  $l$ 
  using XS-nem U(2) by auto
  define uli where  $uli \equiv (\lambda l. from-nat-into\ (U\ l))$ 
  have uli-into:  $uli\ l\ i \in U\ l$  for  $l\ i$ 
  by (simp add: Ul-nem from-nat-into uli-def)
  hence uli-diam:  $mdiameter\ (uli\ l\ i) < 1 / Suc\ l$  for  $l\ i$ 
  using U(3) by auto
  have uli-un:  $M = (\bigcup i. uli\ l\ i)$  for  $l$ 
  by (auto simp: range-from-nat-into[OF Ul-nem[of l] U(1)] uli-def U)
  define Cmn where  $Cmn \equiv (\lambda m\ n. \{f \in mspace\ (cfunspace\ X\ Self).$ 
 $\forall x \in topspace\ X. \forall y \in topspace\ X. dx\ x\ y < 1 / (Suc\ m) \longrightarrow d\ (f\ x)\ (f\ y) < 1$ 
 $/\ Suc\ n\})$ 
  define fmnls where  $fmnls \equiv (\lambda m\ n\ l\ s. SOME\ f. f \in Cmn\ m\ n \wedge (\forall j < km$ 
 $m. f\ (xmk\ m\ j) \in uli\ l\ (s\ j)))$ 
  define Dmnl where  $Dmnl \equiv (\lambda m\ n\ l. \{fmnls\ m\ n\ l\ s \mid s. s \in \{..<km\ m\} \rightarrow_E$ 
 $UNIV \wedge (\exists f \in Cmn\ m\ n. (\forall j < km\ m. f\ (xmk\ m\ j) \in uli\ l\ (s\ j)))\})$ 
  have in-Dmnl:  $fmnls\ m\ n\ l\ s \in Dmnl\ m\ n\ l$  if  $s \in \{..<km\ m\} \rightarrow_E UNIV$   $f \in$ 
 $Cmn\ m\ n \forall j < km\ m. f\ (xmk\ m\ j) \in uli\ l\ (s\ j)$  for  $m\ n\ l\ s\ f$ 
  using Dmnl-def that by blast
  define Dmn where  $Dmn \equiv (\lambda m\ n. \bigcup l. Dmnl\ m\ n\ l)$ 
  have Dmn-subset:  $Dmn\ m\ n \subseteq Cmn\ m\ n$  for  $m\ n$ 
  proof –
  have  $Dmnl\ m\ n\ l \subseteq Cmn\ m\ n$  for  $l$ 
  by (auto simp: Dmnl-def fmnls-def someI[of  $\lambda f. f \in Cmn\ m\ n \wedge (\forall j < km$ 
 $m. f\ (xmk\ m\ j) \in uli\ l\ (-\ j))$ ])
  thus ?thesis by (auto simp: Dmn-def)
  qed
  have c-Dmn:  $countable\ (Dmn\ m\ n)$  for  $m\ n$ 
  proof –
  have  $countable\ (Dmnl\ m\ n\ l)$  for  $l$ 
  proof –
  have  $1 : Dmnl\ m\ n\ l \subseteq (\lambda s. fmnls\ m\ n\ l\ s) \text{ ‘ } (\{..<km\ m\} \rightarrow_E UNIV)$ 
  by (auto simp: Dmnl-def)
  have  $countable\ \dots$ 
  by (auto intro!: countable-PiE)
  with  $1$  show ?thesis
  using countable-subset by blast
  qed
  thus ?thesis
  by (auto simp: Dmn-def)

```

```

qed
have claim:  $\exists g \in Dmn\ m\ n. \forall y \in Xm\ m. d\ (f\ y)\ (g\ y) < e$  if  $f: f \in Cmn\ m\ n$ 
and  $e: e > 0$  for  $f\ m\ n\ e$ 
proof -
  obtain  $l$  where  $l: 1 / Suc\ l < e$ 
  using  $e\ nat\ approx\ posE$  by blast
  define  $s$  where  $s \equiv (\lambda i \in \{.. < km\ m\}. SOME\ j. f\ (xmk\ m\ i) \in uli\ l\ j)$ 
  have  $s1: s \in \{.. < km\ m\} \rightarrow_E\ UNIV$  by (simp add:  $s\ def$ )
  have  $s2: \forall i < km\ m. f\ (xmk\ m\ i) \in uli\ l\ (s\ i)$ 
  proof -
    fix  $i$ 
    have  $f\ (xmk\ m\ i) \in uli\ l\ (SOME\ j. f\ (xmk\ m\ i) \in uli\ l\ j)$  for  $i$ 
    proof (rule someI-ex)
      have  $xmk\ m\ i \in topspace\ X$ 
      using  $Xm(2)\ xmk\ into$  by auto
      hence  $f\ (xmk\ m\ i) \in M$ 
      using  $f$  by (auto simp:  $Cmn\ def\ continuous\ map\ def$ )
      thus  $\exists x. f\ (xmk\ m\ i) \in uli\ l\ x$ 
      using  $uli\ un$  by auto
    qed
    thus ?thesis
    by (auto simp:  $s\ def$ )
  qed
  have  $fmnl: fmnl\ m\ n\ l\ s \in Cmn\ m\ n \wedge (\forall j < km\ m. fmnl\ m\ n\ l\ s\ (xmk\ m\ j) \in uli\ l\ (s\ j))$ 
  by (simp add:  $fmnl\ def, rule\ someI[where\ x=f], auto\ simp: s2\ f$ )
  show  $\exists g \in Dmn\ m\ n. \forall y \in Xm\ m. d\ (f\ y)\ (g\ y) < e$ 
  proof (safe intro!: bestI[where  $x=fmnl\ m\ n\ l\ s$ ])
    fix  $y$ 
    assume  $y: y \in Xm\ m$ 
    then obtain  $i$  where  $i: i < km\ m\ xmk\ m\ i = y$ 
    by (meson  $xmk\ bij[of\ m]\ bij\ betw\ iff\ bijections\ lessThan\ iff$ )
    have  $f\ y \in uli\ l\ (s\ i)\ fmnl\ m\ n\ l\ s\ y \in uli\ l\ (s\ i)$ 
    using  $i(1)\ s2\ fmnl$  by (auto simp:  $i(2)[symmetric]$ )
    moreover have  $f\ y \in M\ fmnl\ m\ n\ l\ s\ y \in M$ 
    using  $f\ fmnl\ y\ Xm(2)[of\ m]$  by (auto simp:  $Cmn\ def\ continuous\ map\ def$ )
    ultimately have  $ennreal\ (d\ (f\ y)\ (fmnl\ m\ n\ l\ s\ y)) \leq mdiameter\ (uli\ l\ (s\ i))$ 
    by (auto intro!:  $mdiameter\ is\ sup$ )
    also have  $\dots < ennreal\ (1 / Suc\ l)$ 
    by (rule  $uli\ diam$ )
    also have  $\dots < ennreal\ e$ 
    using  $l\ e$  by (auto intro!:  $ennreal\ lessI$ )
    finally show  $d\ (f\ y)\ (fmnl\ m\ n\ l\ s\ y) < e$ 
    by (simp add:  $ennreal\ less\ iff[OF\ nonneg]$ )
  qed (use  $in\ Dmnl[OF\ s1\ f\ s2]\ Dmn\ def$  in auto)
qed
show  $separable\ space\ (m\ topology\ of\ (cfunspace\ X\ Self))$ 
unfolding  $separable\ space\ def2\ mdense\ of\ def2$ 

```

```

proof(safe intro!: exI[where  $x = \bigcup n. (\bigcup m. Dmn\ m\ n)$ ])
  fix f and e :: real
  assume h: f ∈ mspace (cfunspace X Self) 0 < e
  then obtain n where n:1 / Suc n < e / 4
    by (metis zero-less-divide-iff zero-less-numeral nat-approx-posE)
  have ∃ m. ∀ l ≥ m. f ∈ Cmn l n
  proof -
    have uniformly-continuous-map dx.Self Self f
    using continuous-eq-uniformly-continuous-map[of dx.Self Self f] h assms(4)
  by(auto simp: mtopology-of-def dx)
    then obtain δ where δ:δ > 0 ∧ x y. x ∈ topspace X ⇒ y ∈ topspace X ⇒
dx x y < δ ⇒ d (f x) (f y) < 1 / (Suc n)
    by(simp only: uniformly-continuous-map-def) (metis dx.mdist-Self
dx.mspace-Self mdist-Self of-nat-0-less-iff zero-less-Suc zero-less-divide-1-iff)
    then obtain m where m:1 / Suc m < δ
    using nat-approx-posE by blast
    have l: l ≥ m ⇒ 1 / Suc l ≤ 1 / Suc m for l
    by (simp add: frac-le)
    show ?thesis
    using δ(2)[OF - - order.strict-trans[OF - order.strict-trans1[OF l m]]] h
  by(auto simp: Cmn-def)
  qed
  then obtain m where m:f ∈ Cmn m n by auto
  obtain g where g:g ∈ Dmn m n ∧ y. y ∈ Xm m ⇒ d (f y) (g y) < e / 4
    by (metis claim[OF m] h(2) zero-less-divide-iff zero-less-numeral)
  have ∃ n m. ∃ g ∈ Dmn m n. mdist (cfunspace X Self) f g < e
  proof(rule exI[where x=n])
    show ∃ m. ∃ g ∈ Dmn m n. mdist (cfunspace X Self) f g < e
    proof(intro exI[where x=m] beXI[OF - g(1)])
      have g-cm:g ∈ mspace (cfunspace X Self)
      using g(1) Dmn-subset[of m n] by(auto simp: Cmn-def)
      have mdist (cfunspace X Self) f g ≤ 3 * e / 4
      proof(rule mdist-cfunspace-le)
        fix x
        assume x:x ∈ topspace X
        obtain y where y:y ∈ Xm m x ∈ dx.mball y (1 / real (Suc m))
          using Xm(3) x by fastforce
        hence ytop:y ∈ topspace X
          by auto
        have mdist Self (f x) (g x) ≤ d (f x) (f y) + d (f y) (g x)
          using x g-cm h(1) ytop by(auto intro!: triangle simp: continu-
ous-map-def)
        also have ... ≤ d (f x) (f y) + d (f y) (g y) + d (g y) (g x)
          using x g-cm h(1) ytop by(auto intro!: triangle simp: continu-
ous-map-def)
        also have ... ≤ e / 4 + e / 4 + e / 4
      proof -
        have dx y: dx x y < 1 / Suc m dx y x < 1 / Suc m
          using y(2) by(auto simp: dx commute)

```

```

    hence  $d(f x) (f y) < 1 / (\text{Suc } n)$   $d(g y) (g x) < 1 / (\text{Suc } n)$ 
      using  $m x y \text{ top } g \text{ Dmn-subset}[of m n]$  by (auto simp: Cmn-def)
    hence  $d(f x) (f y) < e / 4$   $d(g y) (g x) < e / 4$ 
      using  $n$  by auto
    thus ?thesis
      using  $g(2)[OF y(1)]$  by auto
  qed
  finally show  $\text{mdist } Self (f x) (g x) \leq 3 * e / 4$ 
    by simp
  qed(use h in auto)
  also have  $\dots < e$ 
    using  $h$  by auto
  finally show  $\text{mdist } (\text{cfunspace } X \text{ Self}) f g < e$  .
  qed
  qed
  thus  $\exists y \in \bigcup n m. \text{Dmn } m n. \text{mdist } (\text{cfunspace } X \text{ Self}) f y < e$ 
    by auto
  qed(insert Dmn-subset c-Dmn, unfold Cmn-def, blast)+
  qed
  qed
  qed
end

```

```

context Submetric
begin

```

```

lemma separable-sub:
  assumes separable-space mtopology
  shows separable-space sub.mtopology
  using assms unfolding separable-space-iff-second-countable sub.separable-space-iff-second-countable
  by (auto simp: second-countable-subtopology mtopology-submetric)

```

```

end

```

1.3.5 Discrete Distance

```

lemma(in discrete-metric) separable-space-iff: separable-space disc.mtopology  $\longleftrightarrow$ 
countable M
  by (simp add: mtopology-discrete-metric separable-space-discrete-topology)

```

1.3.6 Binary Product Metric Spaces

We define the L^1 -distance. L^1 -distance and L^2 distance (Euclid distance) generate the same topological space.

definition *prod-dist-L1* $\equiv \lambda d1 d2 (x,y) (x',y'). d1 x x' + d2 y y'$

```

context Metric-space12
begin

```


lemma *prod-L1-metric: Metric-space* ($M1 \times M2$) (*prod-dist-L1* $d1$ $d2$)
proof
fix $x\ y\ z$
assume $x \in M1 \times M2\ y \in M1 \times M2\ z \in M1 \times M2$
then show *prod-dist-L1* $d1\ d2\ x\ z \leq \text{prod-dist-L1 } d1\ d2\ x\ y + \text{prod-dist-L1 } d1\ d2\ y\ z$
by(*auto simp: prod-dist-L1-def*) (*metis* $M1.\text{commute}\ M1.\text{triangle''}\ M2.\text{commute}\ M2.\text{triangle''}\ \text{ab-semigroup-add-class.add-ac}(1)\ \text{add.left-commute}\ \text{add-mono}$)
qed(*auto simp: prod-dist-L1-def add-nonneg-eq-0-iff* $M1.\text{commute}\ M2.\text{commute}$)

sublocale *Prod-metric-L1: Metric-space* $M1 \times M2$ *prod-dist-L1* $d1\ d2$
by(*simp add: prod-L1-metric*)

lemma *prod-dist-L1-geq*:
shows $d1\ x\ y \leq \text{prod-dist-L1 } d1\ d2\ (x,x')\ (y,y')$
 $d2\ x'\ y' \leq \text{prod-dist-L1 } d1\ d2\ (x,x')\ (y,y')$
by(*auto simp: prod-dist-L1-def*)

lemma *prod-dist-L1-ball*:
assumes $(x,x') \in \text{Prod-metric-L1.mball } (a,a')\ \varepsilon$
shows $x \in M1.\text{mball } a\ \varepsilon$
and $x' \in M2.\text{mball } a'\ \varepsilon$
using *assms prod-dist-L1-geq order.strict-trans1* **by** *simp-all blast+*

lemma *prod-dist-L1-ball'*:
assumes $z \in \text{Prod-metric-L1.mball } a\ \varepsilon$
shows $\text{fst } z \in M1.\text{mball } (\text{fst } a)\ \varepsilon$
and $\text{snd } z \in M2.\text{mball } (\text{snd } a)\ \varepsilon$
using *prod-dist-L1-ball[of fst z snd z fst a snd a ε]* *assms* **by** *auto*

lemma *prod-dist-L1-ball1'*: $\text{Prod-metric-L1.mball } (a1,a2)\ (\min\ e1\ e2) \subseteq M1.\text{mball } a1\ e1 \times M2.\text{mball } a2\ e2$
using *prod-dist-L1-geq order.strict-trans1* **by** *auto blast+*

lemma *prod-dist-L1-ball1*:
assumes $b1 \in M1.\text{mball } a1\ e1\ b2 \in M2.\text{mball } a2\ e2$
shows $\exists\ e12 > 0. \text{Prod-metric-L1.mball } (b1,b2)\ e12 \subseteq M1.\text{mball } a1\ e1 \times M2.\text{mball } a2\ e2$
proof –
obtain $ea1\ ea2$ **where** $ea1 > 0\ ea2 > 0\ M1.\text{mball } b1\ ea1 \subseteq M1.\text{mball } a1\ e1\ M2.\text{mball } b2\ ea2 \subseteq M2.\text{mball } a2\ e2$
using *assms* **by** (*meson* $M1.\text{openin-mball}\ M1.\text{openin-mtopology}\ M2.\text{openin-mball}\ M2.\text{openin-mtopology}$)
thus *?thesis*
using *prod-dist-L1-ball1'[of b1 b2 ea1 ea2]* **by**(*auto intro!: exI[where $x = \min\ ea1\ ea2$]*)
qed

lemma *prod-dist-L1-ball2'*:
 $M1.mball\ a1\ e1 \times M2.mball\ a2\ e2 \subseteq Prod\text{-}metric\text{-}L1.mball\ (a1, a2)\ (e1 + e2)$
by *auto* (*auto simp: prod-dist-L1-def*)

lemma *prod-dist-L1-ball2*:
assumes $(b1, b2) \in Prod\text{-}metric\text{-}L1.mball\ (a1, a2)\ e12$
shows $\exists e1 > 0. \exists e2 > 0. M1.mball\ b1\ e1 \times M2.mball\ b2\ e2 \subseteq Prod\text{-}metric\text{-}L1.mball\ (a1, a2)\ e12$
proof –
obtain $e12' > 0$ **where** $Prod\text{-}metric\text{-}L1.mball\ (b1, b2)\ e12' \subseteq Prod\text{-}metric\text{-}L1.mball\ (a1, a2)\ e12$
by (*metis assms Prod-metric-L1.openin-mball Prod-metric-L1.openin-mtopology*)
thus *?thesis*
using *prod-dist-L1-ball2'*[*of b1 e12' / 2 b2 e12' / 2*] **by** (*auto intro!: exI*[**where** $x = e12' / 2$])
qed

lemma *prod-dist-L1-mtopology*:
 $Prod\text{-}metric\text{-}L1.mtopology = prod\text{-}topology\ M1.mtopology\ M2.mtopology$
proof –
have $Prod\text{-}metric\text{-}L1.mtopology = topology\text{-}generated\text{-}by\ \{ M1.mball\ a1\ e1 \times M2.mball\ a2\ e2 \mid a1\ a2\ e1\ e2. a1 \in M1 \wedge a2 \in M2 \wedge e1 > 0 \wedge e2 > 0 \}$
unfolding *base-is-subbase*[*OF Prod-metric-L1.mtopology-base-in-balls, simplified subbase-in-def*]
proof (*rule topology-generated-by-eq*)
fix U
assume $U \in \{ M1.mball\ a1\ e1 \times M2.mball\ a2\ e2 \mid a1\ a2\ e1\ e2. a1 \in M1 \wedge a2 \in M2 \wedge 0 < e1 \wedge 0 < e2 \}$
then obtain $a1\ e1\ a2\ e2$ **where** *hae*:
 $U = M1.mball\ a1\ e1 \times M2.mball\ a2\ e2\ a1 \in M1\ a2 \in M2\ 0 < e1\ 0 < e2$
by *auto*
show $openin\ (topology\text{-}generated\text{-}by\ \{ Prod\text{-}metric\text{-}L1.mball\ a\ \varepsilon \mid a\ \varepsilon. a \in M1 \times M2 \wedge 0 < \varepsilon \})\ U$
unfolding *base-is-subbase*[*OF Prod-metric-L1.mtopology-base-in-balls, simplified subbase-in-def, symmetric*] $Prod\text{-}metric\text{-}L1.openin\text{-}mtopology\ hae(1)$
using *prod-dist-L1-ball1*[*of - a1 e1 - a2 e2*] **by** *fastforce*
next
fix U
assume $U \in \{ Prod\text{-}metric\text{-}L1.mball\ a\ \varepsilon \mid a\ \varepsilon. a \in M1 \times M2 \wedge 0 < \varepsilon \}$
then obtain $a1\ a2\ \varepsilon$ **where** *hae*:
 $U = Prod\text{-}metric\text{-}L1.mball\ (a1, a2)\ \varepsilon\ a1 \in M1\ a2 \in M2\ 0 < \varepsilon$
by *auto*
show $openin\ (topology\text{-}generated\text{-}by\ \{ M1.mball\ a1\ e1 \times M2.mball\ a2\ e2 \mid a1\ a2\ e1\ e2. a1 \in M1 \wedge a2 \in M2 \wedge 0 < e1 \wedge 0 < e2 \})\ U$
unfolding *openin-subopen*[*of - Prod-metric-L1.mball (a1, a2) ε*] *hae(1)*
proof *safe*
fix $b1\ b2$
assume $h: (b1, b2) \in Prod\text{-}metric\text{-}L1.mball\ (a1, a2)\ \varepsilon$
from *prod-dist-L1-ball2*[*OF this*] **obtain** $e1\ e2$ **where** $e1 > 0\ e2 > 0$ $M1.mball$

```

b1 e1 × M2.mball b2 e2 ⊆ Prod-metric-L1.mball (a1, a2) ε
  by metis
  with h show ∃ T. openin (topology-generated-by {M1.mball a1 e1 × M2.mball
a2 e2 | a1 a2 e1 e2. a1 ∈ M1 ∧ a2 ∈ M2 ∧ 0 < e1 ∧ 0 < e2}) T ∧ (b1, b2) ∈
T ∧ T ⊆ Prod-metric-L1.mball (a1, a2) ε
    unfolding openin-topology-generated-by-iff
    by(auto intro!: generate-topology-on.Basis exI[where x=M1.mball b1 e1 ×
M2.mball b2 e2])
  qed
qed
also have ... = prod-topology M1.mtopology M2.mtopology
proof –
  have {M1.mball a ε × M2.mball a' ε' | a a' ε ε'. a ∈ M1 ∧ a' ∈ M2 ∧ 0 < ε
∧ 0 < ε'} = {U × V | U V. U ∈ {M1.mball a ε | a ε. a ∈ M1 ∧ 0 < ε} ∧ V ∈
{M2.mball a ε | a ε. a ∈ M2 ∧ 0 < ε}}
    by blast
  thus ?thesis
  unfolding base-is-subbase[OF M1.mtopology-base-in-balls,simplified subbase-in-def]
base-is-subbase[OF M2.mtopology-base-in-balls,simplified subbase-in-def]
  by(simp only: prod-topology-generated-by[symmetric])
qed
finally show ?thesis .
qed

```

lemma *prod-dist-L1-limitin-iff: limitin Prod-metric-L1.mtopology zn z sequentially*
 \longleftrightarrow *limitin M1.mtopology (λn. fst (zn n)) (fst z) sequentially ∧ limitin M2.mtopology*
(λn. snd (zn n)) (snd z) sequentially

```

proof safe
  assume h:limitin Prod-metric-L1.mtopology zn z sequentially
  show limitin M1.mtopology (λn. fst (zn n)) (fst z) sequentially
    limitin M2.mtopology (λn. snd (zn n)) (snd z) sequentially
  unfolding M1.limit-metric-sequentially M2.limit-metric-sequentially
proof safe
  fix e :: real
  assume e: 0 < e
  with h obtain N where N: ∧n. n ≥ N ⇒ zn n ∈ M1 × M2 ∧ n. n ≥ N
 $\implies$  prod-dist-L1 d1 d2 (zn n) z < e
    by(simp only: Prod-metric-L1.limit-metric-sequentially metis)
  show ∃ N. ∀ n ≥ N. fst (zn n) ∈ M1 ∧ d1 (fst (zn n)) (fst z) < e
    ∃ N. ∀ n ≥ N. snd (zn n) ∈ M2 ∧ d2 (snd (zn n)) (snd z) < e
  proof(safe intro!: exI[where x=N])
    fix n
    assume N ≤ n
    from N[OF this]
    show d1 (fst (zn n)) (fst z) < e d2 (snd (zn n)) (snd z) < e
      using order.strict-trans1[OF prod-dist-L1-geq(1)[of fst (zn n) fst z snd (zn
n) snd z]] order.strict-trans1[OF prod-dist-L1-geq(2)[of snd (zn n) snd z fst (zn n)
fst z]]
      by auto

```

```

    qed(use N(1)[simplified mem-Times-iff] in auto)
  qed(use h Prod-metric-L1.limit-metric-sequentially in auto)
next
  assume h:limitin M1.mtopology (λn. fst (zn n)) (fst z) sequentially
    limitin M2.mtopology (λn. snd (zn n)) (snd z) sequentially
  show limitin Prod-metric-L1.mtopology zn z sequentially
    unfolding Prod-metric-L1.limit-metric-sequentially
  proof safe
    fix e :: real
    assume e: 0 < e
    with h obtain N1 N2 where N: ∧n. n ≥ N1 ⇒ fst (zn n) ∈ M1 ∧n. n ≥
N1 ⇒ d1 (fst (zn n)) (fst z) < e / 2
      ∧n. n ≥ N2 ⇒ snd (zn n) ∈ M2 ∧n. n ≥ N2 ⇒ d2 (snd (zn n)) (snd z)
< e / 2
    unfolding M1.limit-metric-sequentially M2.limit-metric-sequentially
    using half-gt-zero by metis
    thus ∃ N. ∀ n ≥ N. zn n ∈ M1 × M2 ∧ prod-dist-L1 d1 d2 (zn n) z < e
    by(fastforce intro!: exI[where x=max N1 N2] simp: prod-dist-L1-def split-beta'
mem-Times-iff)
    qed(auto simp: mem-Times-iff h[simplified M1.limit-metric-sequentially M2.limit-metric-sequentially])
  qed

```

lemma *prod-dist-L1-MCauchy-iff*: $Prod\text{-metric-L1.MCauchy } zn \longleftrightarrow M1.MCauchy$
 $(\lambda n. fst (zn n)) \wedge M2.MCauchy (\lambda n. snd (zn n))$

```

proof safe
  assume h:Prod-metric-L1.MCauchy zn
  show M1.MCauchy (λn. fst (zn n)) M2.MCauchy (λn. snd (zn n))
    unfolding M1.MCauchy-def M2.MCauchy-def
  proof safe
    fix e :: real
    assume 0 < e
    with h obtain N where N: ∧n m. N ≤ n ⇒ N ≤ m ⇒ prod-dist-L1 d1 d2
(zn n) (zn m) < e
    by(fastforce simp: Prod-metric-L1.MCauchy-def)
    show ∃ N. ∀ n n'. N ≤ n ⇒ N ≤ n' ⇒ d1 (fst (zn n)) (fst (zn n')) < e
    ∃ N. ∀ n n'. N ≤ n ⇒ N ≤ n' ⇒ d2 (snd (zn n)) (snd (zn n')) < e
    proof(safe intro!: exI[where x=N])
      fix n m
      assume n ≥ N m ≥ N
      from N[OF this]
      show d1 (fst (zn n)) (fst (zn m)) < e d2 (snd (zn n)) (snd (zn m)) < e
        using order.strict-trans1[OF prod-dist-L1-geq(1)[of fst (zn n) fst (zn m)
snd (zn n) snd (zn m)]] order.strict-trans1[OF prod-dist-L1-geq(2)[of snd (zn n)
snd (zn m) fst (zn n) fst (zn m)]]
        by auto
    qed
  next
  have ∧n. zn n ∈ M1 × M2
    using h by(auto simp: Prod-metric-L1.MCauchy-def)

```

```

    thus fst (zn n) ∈ M1 snd (zn n) ∈ M2 for n
      by (auto simp: mem-Times-iff)
  qed
next
assume h:M1.MCauchy (λn. fst (zn n)) M2.MCauchy (λn. snd (zn n))
show Prod-metric-L1.MCauchy zn
  unfolding Prod-metric-L1.MCauchy-def
proof safe
  fix e :: real
  assume 0 < e
  with h obtain N1 N2 where ∧n m. n ≥ N1 ⇒ m ≥ N1 ⇒ d1 (fst (zn
n)) (fst (zn m)) < e / 2
    ∧n m. n ≥ N2 ⇒ m ≥ N2 ⇒ d2 (snd (zn n)) (snd (zn m)) < e / 2
  unfolding M1.MCauchy-def M2.MCauchy-def using half-gt-zero by metis
  thus ∃N. ∀n n'. N ≤ n → N ≤ n' → prod-dist-L1 d1 d2 (zn n) (zn n') < e
  by (fastforce intro!: exI [where x=max N1 N2] simp: prod-dist-L1-def split-beta')
next
fix x y n
assume 1:(x,y) = zn n
have fst (zn n) ∈ M1 snd (zn n) ∈ M2
  using h[simplified M1.MCauchy-def M2.MCauchy-def] by auto
thus x ∈ M1 y ∈ M2
  by (simp-all add: 1[symmetric])
qed
qed
end

```

1.3.7 Sum Metric Spaces

```

locale Sum-metric =
  fixes I :: 'i set
  and Mi :: 'i ⇒ 'a set
  and di :: 'i ⇒ 'a ⇒ 'a ⇒ real
  assumes Mi-disj: disjoint-family-on Mi I
  and d-nonneg: ∧i x y. 0 ≤ di i x y
  and d-bounded: ∧i x y. di i x y < 1
  and Md-metric: ∧i. i ∈ I ⇒ Metric-space (Mi i) (di i)
begin

abbreviation M ≡ ∪i∈I. Mi i

lemma Mi-inj-on:
  assumes i ∈ I j ∈ I a ∈ Mi i a ∈ Mi j
  shows i = j
  using Mi-disj assms by (auto simp: disjoint-family-on-def)

definition sum-dist :: ['a, 'a] ⇒ real where
sum-dist x y ≡ (if x ∈ M ∧ y ∈ M then (if ∃i∈I. x ∈ Mi i ∧ y ∈ Mi i then di

```

(THE $i. i \in I \wedge x \in Mi\ i \wedge y \in Mi\ i) x\ y\ \text{else}\ 1) \text{ else}\ 0$)

lemma *sum-dist-simps*:

shows $\bigwedge i. \llbracket i \in I; x \in Mi\ i; y \in Mi\ i \rrbracket \implies \text{sum-dist } x\ y = di\ i\ x\ y$
and $\bigwedge i\ j. \llbracket i \in I; j \in I; i \neq j; x \in Mi\ i; y \in Mi\ j \rrbracket \implies \text{sum-dist } x\ y = 1$
and $\bigwedge i. \llbracket i \in I; y \in M; x \in Mi\ i; y \notin Mi\ i \rrbracket \implies \text{sum-dist } x\ y = 1$
and $\bigwedge i. \llbracket i \in I; x \in M; y \in Mi\ i; x \notin Mi\ i \rrbracket \implies \text{sum-dist } x\ y = 1$
and $x \notin M \implies \text{sum-dist } x\ y = 0\ y \notin M \implies \text{sum-dist } x\ y = 0$

proof –

{ **fix** i

assume $h: i \in I\ x \in Mi\ i\ y \in Mi\ i$

then have $\text{sum-dist } x\ y = di\ i\ x\ y$ (THE $i. i \in I \wedge x \in Mi\ i \wedge y \in Mi\ i) x\ y$

by(*auto simp: sum-dist-def*)

also have $\dots = di\ i\ x\ y$

proof –

have (THE $i. i \in I \wedge x \in Mi\ i \wedge y \in Mi\ i) = i$

using *Mi-disj h by(auto intro!: the1-equality simp: disjoint-family-on-def)*

thus *?thesis by simp*

qed

finally show $\text{sum-dist } x\ y = di\ i\ x\ y$. }

show $\bigwedge i\ j. \llbracket i \in I; j \in I; i \neq j; x \in Mi\ i; y \in Mi\ j \rrbracket \implies \text{sum-dist } x\ y = 1$

$\bigwedge i. \llbracket i \in I; y \in M; x \in Mi\ i; y \notin Mi\ i \rrbracket \implies \text{sum-dist } x\ y = 1$ $\bigwedge i. \llbracket i \in I; x \in M; y \in Mi\ i; x \notin Mi\ i \rrbracket \implies \text{sum-dist } x\ y = 1$

$x \notin M \implies \text{sum-dist } x\ y = 0\ y \notin M \implies \text{sum-dist } x\ y = 0$

using *Mi-disj by(auto simp: sum-dist-def disjoint-family-on-def dest: Mi-inj-on)*

qed

lemma *sum-dist-if-less1*:

assumes $i \in I\ x \in Mi\ i\ y \in M\ \text{sum-dist } x\ y < 1$

shows $y \in Mi\ i$

using *assms sum-dist-simps(3) by fastforce*

lemma *inM-cases*:

assumes $x \in M\ y \in M$

and $\bigwedge i. \llbracket i \in I; x \in Mi\ i; y \in Mi\ i \rrbracket \implies P\ x\ y$

and $\bigwedge i\ j. \llbracket i \in I; j \in I; i \neq j; x \in Mi\ i; y \in Mi\ j; x \neq y \rrbracket \implies P\ x\ y$

shows $P\ x\ y$ **using** *assms by auto*

sublocale *Sum-metric: Metric-space M sum-dist*

proof

fix $x\ y$

assume $x \in M\ y \in M$

then show $\text{sum-dist } x\ y = 0 \iff x = y$

by(*rule inM-cases, insert Md-metric*) (*auto simp: sum-dist-simps Metric-space-def*)

next

{ **fix** $i\ x\ y$

assume $h: i \in I\ x \in Mi\ i\ y \in Mi\ i$

then have $\text{sum-dist } x\ y = di\ i\ x\ y$

$\text{sum-dist } y\ x = di\ i\ x\ y$

```

    using Md-metric by(auto simp: sum-dist-simps Metric-space-def) }
  thus  $\bigwedge x y. \text{sum-dist } x y = \text{sum-dist } y x$ 
    by (metis (no-types, lifting) sum-dist-def)
next
show  $1: \bigwedge x y. 0 \leq \text{sum-dist } x y$ 
  using d-nonneg by(simp add: sum-dist-def)
fix  $x y z$ 
assume  $h: x \in M y \in M z \in M$ 
show  $\text{sum-dist } x z \leq \text{sum-dist } x y + \text{sum-dist } y z$  (is ?lhs  $\leq$  ?rhs)
proof(rule inM-cases[OF h(1,3)])
  fix  $i$ 
  assume  $h': i \in I x \in M_i i z \in M_i$ 
  consider  $y \in M_i \mid y \notin M_i$  by auto
  thus  $?lhs \leq ?rhs$ 
  proof cases
    case 1
    with  $h' \text{Md-metric [OF } h'(1)]$  show  $?thesis$ 
      by(simp add: sum-dist-simps Metric-space-def)
  next
    case 2
    with  $h' h(2) \text{d-bounded[of } i x z] 1[\text{of } y z]$ 
    show  $?thesis$ 
      by(auto simp: sum-dist-simps)
  qed
next
fix  $i j$ 
assume  $h': i \in I j \in I i \neq j x \in M_i i z \in M_j$ 
consider  $y \notin M_i \mid y \notin M_j$ 
  using  $h' h(2) \text{Mi-disj}$  by(auto simp: disjoint-family-on-def)
  thus  $?lhs \leq ?rhs$ 
  by (cases, insert 1[of x y] 1[of y z] h' h(2)) (auto simp: sum-dist-simps)
qed
qed

lemma sum-dist-le1: sum-dist x y  $\leq$  1
  using d-bounded[of - x y] by(auto simp: sum-dist-def less-eq-real-def)

lemma sum-dist-ball-eq-ball:
  assumes  $i \in I e \leq 1 x \in M_i$ 
  shows Metric-space.mball (M_i i) (d_i i) x e = Sum-metric.mball x e
proof -
  interpret  $m: \text{Metric-space } M_i i d_i i$ 
  by(simp add: assms Md-metric)
  show  $?thesis$ 
    using assms sum-dist-simps(1)[OF assms(1) assms(3)] sum-dist-if-less1[OF
assms(1,3)]
    by(fastforce simp: Sum-metric.mball-def)
qed

```

```

lemma ball-le-sum-dist-ball:
  assumes  $i \in I$ 
  shows  $\text{Metric-space.mball } (M\ i) (d\ i) x\ e \subseteq \text{Sum-metric.mball } x\ e$ 
proof -
  interpret  $m: \text{Metric-space } M\ i\ d\ i$ 
  by(simp add: assms  $Md\text{-metric}$ )
  show ?thesis
  using assms by(auto simp: sum-dist-simps)
qed

lemma openin-mtopology-iff:
  openin  $\text{Sum-metric.mtopology } U \longleftrightarrow U \subseteq M \wedge (\forall i \in I. \text{openin } (\text{Metric-space.mtopology } (M\ i) (d\ i)) (U \cap M\ i))$ 
proof safe
  fix  $i$ 
  assume  $h: \text{openin } \text{Sum-metric.mtopology } U\ i \in I$ 
  then interpret  $m: \text{Metric-space } M\ i\ d\ i$ 
  using  $Md\text{-metric}$  by simp
  show openin  $m.mtopology (U \cap M\ i)$ 
  unfolding  $m.\text{openin-mtopology}$ 
proof safe
  fix  $x$ 
  assume  $x: x \in U\ x \in M\ i$ 
  with  $h$  obtain  $e$  where  $e: e > 0\ \text{Sum-metric.mball } x\ e \subseteq U$ 
  by(auto simp:  $\text{Sum-metric.openin-mtopology}$ )
  show  $\exists \varepsilon > 0. m.\text{mball } x\ \varepsilon \subseteq U \cap M\ i$ 
  proof(safe intro!: exI[where  $x=e$ ])
    fix  $y$ 
    assume  $y \in m.\text{mball } x\ e$ 
    with  $h(2)$  have  $y \in \text{Sum-metric.mball } x\ e$ 
    by(auto simp: sum-dist-simps)
    with  $e$  show  $y \in U$  by blast
  qed(use  $e$  in auto)
qed
next
  show  $\bigwedge x. \text{openin } \text{Sum-metric.mtopology } U \implies x \in U \implies x \in M$ 
  by(auto simp:  $\text{Sum-metric.openin-mtopology}$ )
next
  assume  $h: U \subseteq M\ \forall i \in I. \text{openin } (\text{Metric-space.mtopology } (M\ i) (d\ i)) (U \cap M\ i)$ 
  show openin  $\text{Sum-metric.mtopology } U$ 
  unfolding  $\text{Sum-metric.openin-mtopology}$ 
proof safe
  fix  $x$ 
  assume  $x: x \in U$ 
  then obtain  $i$  where  $i: i \in I\ x \in M\ i$ 
  using  $h(1)$  by auto
  then interpret  $m: \text{Metric-space } M\ i\ d\ i$ 

```



```

    using Md-metric by simp
  obtain e where e:  $e > 0$  m.mball x e  $\subseteq U \cap Mi\ i$ 
    using i h(2) by (meson IntI m.openin-mtopology x)
  show  $\exists \varepsilon > 0. Sum\text{-metric.mball } x\ \varepsilon \subseteq U$ 
  proof(safe intro!: exI[where x=min e 1])
    fix y
    assume y:y  $\in Sum\text{-metric.mball } x\ (min\ e\ 1)$ 
    then show y  $\in U$ 
      using sum-dist-ball-eq-ball[OF i(1) - i(2), of min e 1] e by fastforce
    qed(simp add: e)
  qed(use h(1) in auto)
qed

```

corollary *openin-mtopology-Mi*:

```

  assumes i  $\in I$ 
  shows openin Sum-metric.mtopology (Mi i)
  unfolding openin-mtopology-iff
proof safe
  fix j
  assume j:j  $\in I$ 
  then interpret m: Metric-space Mi j di j
    by(simp add: Md-metric)
  show openin m.mtopology (Mi i  $\cap$  Mi j)
    by (cases i = j, insert assms Mi-disj j) (auto simp: disjoint-family-on-def)
  qed(use assms in auto)

```

corollary *subtopology-mtopology-Mi*:

```

  assumes i  $\in I$ 
  shows subtopology Sum-metric.mtopology (Mi i) = Metric-space.mtopology (Mi
i) (di i)
proof –
  interpret Mi: Metric-space Mi i di i
    by (simp add: Md-metric assms)
  show ?thesis
    unfolding topology-eq openin-subtopology
  proof safe
    fix T
    assume openin Sum-metric.mtopology T
    thus openin Mi.mtopology (T  $\cap$  Mi i)
      using assms by(auto simp: openin-mtopology-iff)
  next
    fix S
    assume h:openin Mi.mtopology S
    show  $\exists T. openin\ Sum\text{-metric.mtopology } T \wedge S = T \cap Mi\ i$ 
    proof(safe intro!: exI[where x=S])
      show openin Sum-metric.mtopology S
        unfolding openin-mtopology-iff
      proof safe
        fix j

```

```

assume  $j:j \in I$ 
then interpret  $Mj$ : Metric-space  $Mi\ j\ di\ j$ 
  using Md-metric by auto
have  $i \neq j \implies S \cap Mi\ j = \{\}$ 
  using openin-subset[OF h] Mi-disj j assms
  by(auto simp: disjoint-family-on-def)
thus openin  $Mj.mtopology\ (S \cap Mi\ j)$ 
  by(cases i = j) (use openin-subset[OF h] h in auto)
qed(use openin-subset[OF h] assms in auto)
qed(use openin-subset[OF h] assms in auto)
qed
qed

```

lemma *limitin-Mi-limitin-M*:

```

assumes  $i \in I$  limitin (Metric-space.mtopology ( $Mi\ i$ ) ( $di\ i$ ))  $xn\ x$  sequentially
shows limitin Sum-metric.mtopology  $xn\ x$  sequentially
proof –
interpret  $m$ : Metric-space  $Mi\ i\ di\ i$ 
  by(simp add: assms Md-metric)
  {
    fix  $e :: real$ 
    assume  $e > 0$ 
    then obtain  $N$  where  $\bigwedge n. n \geq N \implies xn\ n \in m.mball\ x\ e$ 
      using assms(2) m.commute m.limit-metric-sequentially by fastforce
    hence  $\exists N. \forall n \geq N. xn\ n \in Sum-metric.mball\ x\ e$ 
      using ball-le-sum-dist-ball[OF assms(1),of x e]
      by(fastforce intro!: exI[where  $x=N$ ]) }
    thus ?thesis
  }
by (metis Sum-metric.commute Sum-metric.in-mball Sum-metric.limit-metric-sequentially
UN-I m.limitin-mspace assms)
qed

```

lemma *limitin-M-limitin-Mi*:

```

assumes limitin Sum-metric.mtopology  $xn\ x$  sequentially
shows  $\exists i \in I. \textit{limitin}$  (Metric-space.mtopology ( $Mi\ i$ ) ( $di\ i$ ))  $xn\ x$  sequentially
proof –
obtain  $i$  where  $i \in I\ x \in Mi\ i$ 
  using assms by (meson Sum-metric.limitin-mspace UN-E)
then interpret  $m$ : Metric-space  $Mi\ i\ di\ i$ 
  by(simp add: Md-metric)
obtain  $N$  where  $N: \bigwedge n. n \geq N \implies sum-dist\ (xn\ n)\ x < 1\ \bigwedge n. n \geq N \implies$ 
 $(xn\ n) \in M$ 
  using assms by (metis d-bounded i(2) m.mdist-zero Sum-metric.limit-metric-sequentially)
hence  $xni: n \geq N \implies xn\ n \in Mi\ i$  for  $n$ 
  using assms by(auto intro!: sum-dist-if-less1[OF i,of xn n] simp: Sum-metric.commute)
show ?thesis
proof(safe intro!: bexI[where  $x=i$ ]  $i$ )
  show limitin  $m.mtopology\ xn\ x$  sequentially
    unfolding m.limit-metric-sequentially

```

```

proof safe
  fix e :: real
  assume e: 0 < e
  then obtain N' where N':  $\bigwedge n. n \geq N' \implies \text{sum-dist } (xn \ n) \ x < e$ 
    using assms by (meson Sum-metric.limit-metric-sequentially)
  hence  $n \geq \max N \ N' \implies di \ i \ (xn \ n) \ x < e$  for n
    using sum-dist-simps(1)[OF i(1) xni[of n] i(2),symmetric] by auto
  thus  $\exists N. \forall n \geq N. xn \ n \in Mi \ i \wedge di \ i \ (xn \ n) \ x < e$ 
    using xni by(auto intro!: exI[where x=max N N'])
qed fact
qed
qed

```

```

lemma MCauchy-Mi-MCauchy-M:
  assumes  $i \in I$  Metric-space.MCauchy (Mi i) (di i) xn
  shows Sum-metric.MCauchy xn
proof -
  interpret m: Metric-space Mi i di i
    by(simp add: assms Md-metric)
  have [simp]: $\text{sum-dist } (xn \ n) \ (xn \ m) = di \ i \ (xn \ n) \ (xn \ m)$  for n m
    using assms sum-dist-simps(1)[OF assms(1),of xn n xn m]
    by(auto simp: m.MCauchy-def)
  show ?thesis
    using assms by(auto simp: m.MCauchy-def Sum-metric.MCauchy-def)
qed

```

```

lemma MCauchy-M-MCauchy-Mi:
  assumes Sum-metric.MCauchy xn
  shows  $\exists m. \exists i \in I. \text{Metric-space.MCauchy } (Mi \ i) \ (di \ i) \ (\lambda n. xn \ (n + m))$ 
proof -
  obtain N where N:  $\bigwedge n \ m. n \geq N \implies m \geq N \implies \text{sum-dist } (xn \ n) \ (xn \ m) < 1$ 
    using assms by(fastforce simp: Sum-metric.MCauchy-def)
  obtain i where i:  $i \in I \ \text{and} \ N \in Mi \ i$ 
    by (metis assms Sum-metric.MCauchy-def UNIV-I UN-E image-eqI subsetD)
  then have  $xn: \bigwedge n. n \geq N \implies xn \ n \in Mi \ i$ 
    by (metis N Sum-metric.MCauchy-def assms dual-order.refl range-subsetD sum-dist-if-less1)
  interpret m: Metric-space Mi i di i
    using i Md-metric by auto
  show ?thesis
proof(safe intro!: exI[where x=N] bexI[where x=i])
  show  $m.MCauchy \ (\lambda n. xn \ (n + N))$ 
    unfolding m.MCauchy-def
proof safe
  show 1:  $\bigwedge n. xn \ (n + N) \in Mi \ i$ 
    by(auto intro!: xn)
  fix e :: real
  assume 0 < e
  then obtain N' where N':  $\bigwedge n \ m. n \geq N' \implies m \geq N' \implies \text{sum-dist } (xn \ n)$ 

```

```

(xn m) < e
  using Sum-metric.MCauchy-def assms by blast
  thus  $\exists N'. \forall n n'. N' \leq n \longrightarrow N' \leq n' \longrightarrow di\ i\ (xn\ (n + N))\ (xn\ (n' + N))$ 
  < e
    by(auto intro!: exI[where x=N'] simp: sum-dist-simps(1)[OF i(1) xn
  xn,symmetric])
  qed
  qed fact
  qed

lemma separable-Mi-separable-M:
  assumes countable I  $\wedge i. i \in I \implies separable\ space\ (Metric\ space.mtopology\ (Mi\ i)\ (di\ i))$ 
  shows separable-space Sum-metric.mtopology
  proof -
    obtain Ui where Ui:  $\wedge i. i \in I \implies countable\ (Ui\ i) \wedge i. i \in I \implies dense\ in$ 
    (Metric-space.mtopology (Mi i) (di i)) (Ui i)
    using assms by(simp only: separable-space-def2) metis
    define U where  $U \equiv \bigcup_{i \in I}. Ui\ i$ 
    show separable-space Sum-metric.mtopology
    unfolding separable-space-def2
    proof(safe intro!: exI[where x=U])
      show countable U
      using Ui(1) assms by(auto simp: U-def)
    next
      show Sum-metric.mdense U
      unfolding Sum-metric.mdense-def U-def
    proof safe
      fix i x
      assume  $i \in I\ x \in Ui\ i$ 
      then show  $x \in M$ 
      using Ui(2) by(auto intro!: bexI[where x=i] simp: Md-metric Metric-space.mdense-def2)
    next
      fix i x e
      assume  $h: i \in I\ x \in Mi\ i\ (0 :: real) < e\ Sum\ metric.mball\ x\ e \cap \bigcup (Ui\ 'I)$ 
      = {}
      then interpret m: Metric-space Mi i di i
      by(simp add: Md-metric)
      have  $m.mball\ x\ e \cap Ui\ i \neq \{\}$ 
      using Ui(2)[OF h(1)] h by(auto simp: m.mdense-def)
      hence  $m.mball\ x\ e \cap \bigcup (Ui\ 'I) \neq \{\}$ 
      using h(1) by blast
      thus False
      using ball-le-sum-dist-ball[OF <i ∈ I>,of x e] h(4) by blast
    qed
  qed
  qed
  qed

```

lemma *separable-M-separable-Mi*:
assumes *separable-space Sum-metric.mtopology* $\bigwedge i. i \in I$
shows *separable-space (Metric-space.mtopology (Mi i) (di i))*
using *subtopology-mtopology-Mi[OF assms(2)] separable-space-open-subset[OF*
assms(1) openin-mtopology-Mi[OF assms(2)]]
by *simp*

lemma *mcomplete-Mi-mcomplete-M*:
assumes $\bigwedge i. i \in I \implies \text{Metric-space.mcomplete } (Mi\ i)\ (di\ i)$
shows *Sum-metric.mcomplete*
unfolding *Sum-metric.mcomplete-def*
proof *safe*
fix *xn*
assume *Sum-metric.MCauchy xn*
from *MCauchy-M-MCauchy-Mi[OF this]* **obtain** *m i* **where** *mi*:
 $i \in I \text{ Metric-space.MCauchy } (Mi\ i)\ (di\ i)\ (\lambda n. xn\ (n + m))$
by *metis*
then interpret *m: Metric-space Mi i di i*
by(*simp add: Md-metric*)
from *assms[OF mi(1)] mi(2)* **obtain** *x* **where** *x: limitin m.mtopology* $(\lambda n. xn\ (n + m))\ x\ \text{sequentially}$
by(*auto simp: m.mcomplete-def*)
from *limitin-Mi-limitin-M[OF mi(1) limitin-sequentially-offset-rev[OF this]]*
show $\exists x. \text{limitin } \text{Sum-metric.mtopology } xn\ x\ \text{sequentially}$
by *auto*
qed

lemma *mcomplete-M-mcomplete-Mi*:
assumes *Sum-metric.mcomplete* $i \in I$
shows *Metric-space.mcomplete (Mi i) (di i)*
proof –
interpret *Mi: Metric-space Mi i di i*
using *assms(2) Md-metric* **by** *fastforce*
show *?thesis*
unfolding *Mi.mcomplete-def*
proof *safe*
fix *xn*
assume *xn:Mi.MCauchy xn*
from *MCauchy-Mi-MCauchy-M[OF assms(2) this]* *assms(1)* **obtain** *x* **where**
limitin Sum-metric.mtopology xn x sequentially
by(*auto simp: Sum-metric.mcomplete-def*)
from *limitin-M-limitin-Mi[OF this]* **obtain** *j* **where** $j:j \in I \text{ limitin } (\text{Metric-space.mtopology } (Mi\ j)\ (di\ j))\ xn\ x\ \text{sequentially}$
by *auto*
have $j = i$
proof –
obtain *N* **where** $\bigwedge n. n \geq N \implies xn\ n \in Mi\ j$
by (*metis Md-metric Metric-space.limitin-metric-dist-null eventually-sequentially*
j)

```

    hence  $xn N \in Mi i \cap Mi j$ 
    using  $xn$  by(auto simp: Mi.MCauchy-def)
    with Mi-disj  $j(1)$  assms(2) show ?thesis
    by(fastforce simp: disjoint-family-on-def)
  qed
  thus  $\exists x. \text{limitin } Mi.mtopology \text{ } xn \text{ } x \text{ sequentially}$ 
  using  $j(2)$  by(auto intro!: exI[where  $x=x$ ])
  qed
  qed
end

```

```

lemma sum-metricI:
  fixes  $Si$ 
  assumes disjoint-family-on  $Si I$ 
    and  $\bigwedge i x y. i \notin I \implies 0 \leq di i x y$ 
    and  $\bigwedge i x y. di i x y < 1$ 
    and  $\bigwedge i. i \in I \implies \text{Metric-space } (Si i) (di i)$ 
  shows Sum-metric  $I Si di$ 
  using assms by (metis Metric-space.nonneg Sum-metric-def)

```

end

1.3.8 Product Metric Spaces

```

theory Set-Based-Metric-Product
  imports Set-Based-Metric-Space
begin

```

```

lemma nsum-of-r':
  fixes  $r :: real$ 
  assumes  $r:0 < r r < 1$ 
  shows  $(\sum n. r^{(n+k)} * K) = r^k / (1 - r) * K$ 
  (is ?lhs = -)
proof -
  have ?lhs =  $(\sum n. r^n * K) - (\sum n \in \{..<k\}. r^n * K)$ 
    using  $r$  by(auto intro!: suminf-minus-initial-segment summable-mult2)
  also have  $\dots = 1 / (1 - r) * K - (1 - r^k) / (1 - r) * K$ 
  proof -
    have  $(\sum n \in \{..<k\}. r^n * K) = (1 - r^k) / (1 - r) * K$ 
      using one-diff-power-eq[of  $r k$ ] scale-sum-left[of  $\lambda n. r^n \{..<k\} K, \text{symmetric}$ ]
      by auto
    thus ?thesis
      using  $r$  by(auto simp add: suminf-geometric[of  $r$ ] suminf-mult2[where
 $c=K, \text{symmetric}$ ])
  qed
  finally show ?thesis

```

using r by (*simp add: diff-divide-distrib left-diff-distrib*)
 qed

lemma *nsum-of-r-leq*:

fixes $r :: \text{real}$ and $a :: \text{nat} \Rightarrow \text{real}$
 assumes $r: 0 < r \ r < 1$
 and $a: \bigwedge n. 0 \leq a \ n \ \bigwedge n. a \ n \leq K$
 shows $0 \leq (\sum n. r^{n+k} * a \ (n+l)) \ (\sum n. r^{n+k} * a \ (n+l)) \leq r^k$
 $/ \ (1 - r) * K$
proof –
 have [*simp*]: *summable* $(\lambda n. r^{n+k} * a \ (n+l))$
 apply (*rule summable-comparison-test'*[of $\lambda n. r^{n+k} * K$])
 using $r \ a$ by (*auto intro!: summable-mult2*)
 show $0 \leq (\sum n. r^{n+k} * a \ (n+l))$
 using $r \ a$ by (*auto intro!: suminf-nonneg*)
 have $(\sum n. r^{n+k} * a \ (n+l)) \leq (\sum n. r^{n+k} * K)$
 using $a \ r$ by (*auto intro!: suminf-le summable-mult2*)
 also have $\dots = r^k / (1 - r) * K$
 by (*rule nsum-of-r'*[*OF r*])
 finally show $(\sum n. r^{n+k} * a \ (n+l)) \leq r^k / (1 - r) * K$.
 qed

lemma *nsum-of-r-le*:

fixes $r :: \text{real}$ and $a :: \text{nat} \Rightarrow \text{real}$
 assumes $r: 0 < r \ r < 1$
 and $a: \bigwedge n. 0 \leq a \ n \ \bigwedge n. a \ n \leq K \ \exists n' \geq l. a \ n' < K$
 shows $(\sum n. r^{n+k} * a \ (n+l)) < r^k / (1 - r) * K$
proof –
 obtain n' where $hn': a \ (n' + l) < K$
 using $a(3)$ by (*metis add.commute le-iff-add*)
 define a' where $a' = (\lambda n. \text{if } n = n' + l \text{ then } K \text{ else } a \ n)$
 have $a': \bigwedge n. 0 \leq a' \ n \ \bigwedge n. a' \ n \leq K$
 using $a(1,2)$ *le-trans order.trans*[*OF a(1,2)*][of 0] by (*auto simp: a'-def*)
 have [*simp*]: *summable* $(\lambda n. r^{n+k} * a \ (n+l))$
 apply (*rule summable-comparison-test'*[of $\lambda n. r^{n+k} * K$])
 using $r \ a$ by (*auto intro!: summable-mult2*)
 have [*simp*]: *summable* $(\lambda n. r^{n+k} * a' \ (n+l))$
 apply (*rule summable-comparison-test'*[of $\lambda n. r^{n+k} * K$])
 using $r \ a'$ by (*auto intro!: summable-mult2*)
 have $(\sum n. r^{n+k} * a \ (n+l)) = (\sum n. r^{n+k} * a' \ (n+l)) + (\sum i < n'. r^{i+k} * a \ (i+l))$
 by (*rule suminf-split-initial-segment*) *simp*
 also have $\dots = (\sum n. r^{n+k} * a' \ (n+l)) + (\sum i < n'. r^{i+k} * a \ (i+l)) + r^{n'+k} * a \ (n'+l)$
 by *simp*
 also have $\dots < (\sum n. r^{n+k} * a' \ (n+l)) + (\sum i < n'. r^{i+k} * a \ (i+l)) + r^{n'+k} * a \ (n'+l)$
 using $r \ hn'$ by *auto*
 also have $\dots = (\sum n. r^{n+k} * a' \ (n+l)) + (\sum i < n'. r^{i+k} * a \ (i+l)) + r^{n'+k} * a \ (n'+l)$

$n'. r \hat{\sim}(i + k) * a' (i + l)$
by(*auto simp: a'-def*)
also have ... = $(\sum n. r \hat{\sim}(n + k) * a' (n + l))$
by(*rule suminf-split-initial-segment[symmetric]*) *simp*
also have ... $\leq r \hat{\sim}k / (1 - r) * K$
by(*rule nsum-of-r-leq[OF r a']*)
finally show ?thesis .
qed

definition *product-dist'* :: [*real, 'i set, nat \Rightarrow 'i, 'i \Rightarrow 'a set, 'i \Rightarrow 'a \Rightarrow 'a \Rightarrow real*]
 \Rightarrow (*'i \Rightarrow 'a*) \Rightarrow (*'i \Rightarrow 'a*) \Rightarrow *real* **where**
*product-dist-def: product-dist' r I g Mi di \equiv ($\lambda x y. \text{if } x \in (\prod_E i \in I. Mi i) \wedge y \in (\prod_E i \in I. Mi i) \text{ then } (\sum n. \text{if } g n \in I \text{ then } r \hat{\sim}n * di (g n) (x (g n)) (y (g n)) \text{ else } 0) \text{ else } 0$)*

$$d(x, y) = \sum_{n \in \mathbb{N}} r^n * d_{g_I(i)}(x_{g_I(i)}, y_{g_I(i)}).$$

locale *Product-metric* =
fixes *r :: real*
and *I :: 'i set*
and *f :: 'i \Rightarrow nat*
and *g :: nat \Rightarrow 'i*
and *Mi :: 'i \Rightarrow 'a set*
and *di :: 'i \Rightarrow 'a \Rightarrow 'a \Rightarrow real*
and *K :: real*
assumes *r: 0 < r r < 1*
and *I: countable I*
and *gf-comp-id: $\bigwedge i. i \in I \implies g (f i) = i$*
and *gf-if-finite: finite I \implies bij-betw f I {..*card I*}*
*finite I \implies bij-betw g {..*card I*} I*
and *gf-if-infinite: infinite I \implies bij-betw f I UNIV*
infinite I \implies bij-betw g UNIV I
 $\bigwedge n. infinite I \implies f (g n) = n$
and *Md-metric: $\bigwedge i. i \in I \implies \text{Metric-space } (Mi i) (di i)$*
and *di-nonneg: $\bigwedge i x y. 0 \leq di i x y$*
and *di-bounded: $\bigwedge i x y. di i x y \leq K$*
and *K-pos: 0 < K*

lemma *from-nat-into-to-nat-on-product-metric-pair:*
assumes *countable I*
shows $\bigwedge i. i \in I \implies \text{from-nat-into } I \text{ (to-nat-on } I i) = i$
and *finite I \implies bij-betw (to-nat-on I) I {..*card I*}*
and *finite I \implies bij-betw (from-nat-into I) {..*card I*} I*
and *infinite I \implies bij-betw (to-nat-on I) I UNIV*
and *infinite I \implies bij-betw (from-nat-into I) UNIV I*
and $\bigwedge n. infinite I \implies \text{to-nat-on } I \text{ (from-nat-into } I n) = n$
by(*simp-all add: assms to-nat-on-finite bij-betw-from-nat-into-finite to-nat-on-infinite bij-betw-from-nat-into*)

lemma *product-metric-pair-finite-nat:*

bij-betw id $\{..n\}$ $\{..< \text{card } \{..n\}\}$ *bij-betw id* $\{..< \text{card } \{..n\}\}$ $\{..n\}$
by(*auto simp: bij-betw-def*)

lemma *product-metric-pair-finite-nat'*:

bij-betw id $\{..<n\}$ $\{..< \text{card } \{..<n\}\}$ *bij-betw id* $\{..< \text{card } \{..<n\}\}$ $\{..<n\}$
by(*auto simp: bij-betw-def*)

context *Product-metric*
begin

abbreviation *product-dist* \equiv *product-dist' r I g Mi di*

lemma *nsum-of-rK*: $(\sum n. r^{(n+k)*K}) = r^k / (1 - r) * K$
by(*rule nsum-of-r'[OF r]*)

lemma *i-min*:

assumes $i \in I$ $g\ n = i$
shows $f\ i \leq n$
proof(*cases finite I*)
case $h: \text{True}$
show *?thesis*
proof(*rule ccontr*)
assume $\neg f\ i \leq n$
then have $h0: n < f\ i$ **by** *simp*
have $f\ i \in \{..<\text{card } I\}$
using *bij-betwE[OF gf-if-finite(1)[OF h]] assms(1)* **by** *simp*
moreover have $n \in \{..<\text{card } I\}$ $n \neq f\ i$
using $h0 \langle f\ i \in \{..<\text{card } I\} \rangle$ **by** *auto*
ultimately have $g\ n \neq g\ (f\ i)$
using *bij-betw-imp-inj-on[OF gf-if-finite(2)[OF h]]*
by (*simp add: inj-on-contrad*)
thus *False*
by(*simp add: gf-comp-id[OF assms(1)] assms(2)*)
qed
next
show *infinite I* $\implies f\ i \leq n$
using *assms(2) gf-if-infinite(3)[of n]* **by** *simp*
qed

lemma *g-surj*:

assumes $i \in I$
shows $\exists n. g\ n = i$
using *gf-comp-id[of i] assms* **by** *auto*

lemma *product-dist-summable'[simp]*:

summable $(\lambda n. r^n * di\ (g\ n)\ (x\ (g\ n))\ (y\ (g\ n)))$
apply(*rule summable-comparison-test'[of $\lambda n. r^n * K$]*)
using *r di-nonneg di-bounded K-pos* **by**(*auto intro!: summable-mult2*)

lemma *product-dist-summable*[simp]:
summable ($\lambda n. \text{if } g \ n \in I \text{ then } r \hat{\ } n * di \ (g \ n) \ (x \ (g \ n)) \ (y \ (g \ n)) \ \text{else } 0$)
by(rule *summable-comparison-test'*[of $\lambda n. r \hat{\ } n * di \ (g \ n) \ (x \ (g \ n)) \ (y \ (g \ n))$])
(use r di-nonneg di-bounded K-pos in auto)

lemma *summable-rK*[simp]: *summable* ($\lambda n. r \hat{\ } n * K$)
using *r* **by**(*auto intro!*: *summable-mult2*)

lemma *Product-metric: Metric-space* ($\Pi_E \ i \in I. \ Mi \ i$) *product-dist*
proof –
have h' : $\bigwedge i \ xi \ yi. \ i \in I \implies xi \in Mi \ i \implies yi \in Mi \ i \implies xi = yi \iff di \ i \ xi \ yi = 0$
 $\bigwedge i \ xi \ yi. \ i \in I \implies di \ i \ xi \ yi = di \ i \ yi \ xi$
 $\bigwedge i \ xi \ yi \ zi. \ i \in I \implies xi \in Mi \ i \implies yi \in Mi \ i \implies zi \in Mi \ i \implies di \ i \ xi \ zi \leq di \ i \ xi \ yi + di \ i \ yi \ zi$
using *Md-metric* **by**(*auto simp: Metric-space-def*)
show *?thesis*
proof
show $\bigwedge x \ y. \ 0 \leq \text{product-dist } x \ y$
using *di-nonneg r* **by**(*auto simp: product-dist-def intro! suminf-nonneg product-dist-summable*)
next
fix $x \ y$
assume hxy : $x \in (\Pi_E \ i \in I. \ Mi \ i) \ y \in (\Pi_E \ i \in I. \ Mi \ i)$
show (*product-dist* $x \ y = 0$) $\iff (x = y)$
proof
assume heq : $x = y$
then have (*if* $g \ n \in I \text{ then } r \hat{\ } n * di \ (g \ n) \ (x \ (g \ n)) \ (y \ (g \ n)) \ \text{else } 0$) = 0
for n
using $hxy \ h'(1)$ [of $g \ n \ x \ (g \ n) \ y \ (g \ n)$] **by**(*auto simp: product-dist-def*)
thus *product-dist* $x \ y = 0$
by(*auto simp: product-dist-def*)
next
assume $h0$: *product-dist* $x \ y = 0$
have ($\sum n. \ \text{if } g \ n \in I \text{ then } r \hat{\ } n * di \ (g \ n) \ (x \ (g \ n)) \ (y \ (g \ n)) \ \text{else } 0$) = 0
 $\iff (\forall n. \ (\text{if } g \ n \in I \text{ then } r \hat{\ } n * di \ (g \ n) \ (x \ (g \ n)) \ (y \ (g \ n)) \ \text{else } 0) = 0)$
apply(rule *suminf-eq-zero-iff*)
using *di-nonneg r* **by**(*auto simp: product-dist-def intro! product-dist-summable*)
hence $hn0$: $\bigwedge n. \ (\text{if } g \ n \in I \text{ then } r \hat{\ } n * di \ (g \ n) \ (x \ (g \ n)) \ (y \ (g \ n)) \ \text{else } 0) = 0$
using $h0 \ hxy$ **by**(*auto simp: product-dist-def*)
show $x = y$
proof
fix i
consider $i \in I \mid i \notin I$ **by** *auto*
thus $x \ i = y \ i$
proof *cases*
case 1
from *g-surj*[*OF this*] **obtain** n **where**

```

      hn: g n = i by auto
    have di (g n) (x (g n)) (y (g n)) = 0
      using hn h'(1)[OF 1,of x i y i] hxy hn0[of n] 1 r by simp
    thus ?thesis
      using hn h'(1)[OF 1,of x i y i] hxy 1 by auto
  next
  case 2
  then show ?thesis
    by(simp add: PiE-arb[OF hxy(1) 2] hxy PiE-arb[OF hxy(2) 2])
  qed
qed
qed
next
show product-dist x y = product-dist y x for x y
  using h'(2) by(auto simp: product-dist-def) (metis (no-types, opaque-lifting))
next
fix x y z
assume hxyz:x ∈ (ΠE i∈I. Mi i) y ∈ (ΠE i∈I. Mi i) z ∈ (ΠE i∈I. Mi i)
have (if g n ∈ I then r ^ n * di (g n) (x (g n)) (z (g n)) else 0)
  ≤ (if g n ∈ I then r ^ n * di (g n) (x (g n)) (y (g n)) else 0) + (if g n ∈
I then r ^ n * di (g n) (y (g n)) (z (g n)) else 0) for n
  using h'(3)[of g n x (g n) y (g n) z (g n)] hxyz r
  by(auto simp: distrib-left[of r ^ n,symmetric])
  thus product-dist x z ≤ product-dist x y + product-dist y z
  by(auto simp add: product-dist-def suminf-add[OF product-dist-summable[of x
y] product-dist-summable[of y z]] hxyz intro!: suminf-le summable-add)
  qed
qed

sublocale Product-metric: Metric-space ΠE i∈I. Mi i product-dist
  by(rule Product-metric)

lemma product-dist-leqr: product-dist x y ≤ 1 / (1 - r) * K
proof -
  have product-dist x y ≤ (∑ n. if g n ∈ I then r ^ n * di (g n) (x (g n)) (y (g n))
else 0)
  proof -
    consider x ∈ (ΠE i∈I. Mi i) ∧ y ∈ (ΠE i∈I. Mi i) | ¬ (x ∈ (ΠE i∈I. Mi i)
∧ y ∈ (ΠE i∈I. Mi i)) by auto
    then show ?thesis
    proof cases
      case 1
      then show ?thesis by(auto simp: product-dist-def)
    next
      case 2
      then have product-dist x y = 0
        by(auto simp: product-dist-def)
      also have ... ≤ (∑ n. if g n ∈ I then r ^ n * di (g n) (x (g n)) (y (g n)) else
0)

```

```

    using di-nonneg r by(auto intro!: suminf-nonneg product-dist-summable)
    finally show ?thesis .
  qed
  qed
  also have ... ≤ (∑ n. r^n * di (g n) (x (g n)) (y (g n)))
    using r di-nonneg di-bounded by(auto intro!: suminf-le)
  also have ... ≤ (∑ n. r^n * K)
    using r di-bounded di-nonneg by(auto intro!: suminf-le)
  also have ... = 1 / (1 - r) * K
    using r nsum-of-rK[of 0] by simp
  finally show ?thesis .
  qed

lemma product-dist-geq:
  assumes i ∈ I and g n = i x ∈ (ΠE i∈I. M i) y ∈ (ΠE i∈I. M i)
  shows di i (x i) (y i) ≤ (1/r)^n * product-dist x y
    (is ?lhs ≤ ?rhs)
  proof -
    interpret mi: Metric-space M i di i
    by(rule Md-metric[OF assms(1)])
    have (λm. if m = f i then di (g m) (x (g m)) (y (g m)) else 0) sums di (g (f i))
      (x (g (f i))) (y (g (f i)))
    by(rule sums-single)
    also have ... = ?lhs
    by(simp add: gf-comp-id[OF assms(1)])
    finally have 1:summable (λm. if m = f i then di (g m) (x (g m)) (y (g m)) else
      0)
      ?lhs = (∑ m. (if m = f i then di (g m) (x (g m)) (y (g m)) else 0))
    by(auto simp: sums-iff)
    note 1(2)
    also have ... ≤ (∑ m. (1/r)^n * (if g m ∈ I then r^m * di (g m) (x (g m)) (y
      (g m)) else 0))
    proof(rule suminf-le)
      show summable (λm. (1/r)^n * (if g m ∈ I then r^m * di (g m) (x (g m)) (y
        (g m)) else 0))
      by(auto intro!: product-dist-summable)
    next
      fix k
      have **:1 ≤ (1/r)^n * r^k
      proof -
        have (1/r)^n * r^k = (1/r)^(n-k) * (1/r)^k * r^k
        using r by(simp add: power-diff[OF -i-min[OF assms(1,2)],of 1/r,simplified])
        also have ... = (1/r)^(n-k)
        using r by (simp add: power-one-over)
        finally show ?thesis
        using r by auto
      qed
    next
      have *:g k ∈ I if k = f i
      using gf-comp-id[OF assms(1)] assms(1) that by auto
  
```

```

show (if k = f i then di (g k) (x (g k)) (y (g k)) else 0) ≤ (1/r) ^ n * (if g k
∈ I then r ^ k * di (g k) (x (g k)) (y (g k)) else 0)
using * di-nonneg r ** mult-right-mono[OF **] by(auto simp: vector-space-over-itself.scale-scale[of
(1 / r) ^ n])
qed(simp add: 1)
also have ... = ?rhs
unfolding product-dist-def
using assms by(auto intro!: suminf-mult product-dist-summable)
finally show ?thesis .
qed

```

lemma *limitin-M-iff-limitin-Mi*:

```

shows limitin Product-metric.mtopology xn x sequentially ↔ (∃ N. ∀ n ≥ N.
(∀ i ∈ I. xn n i ∈ Mi i) ∧ (∀ i. i ∉ I → xn n i = undefined)) ∧ (∀ i ∈ I. limitin
(Metric-space.mtopology (Mi i) (di i)) (λ n. xn n i) (x i) sequentially) ∧ x ∈ (ΠE
i ∈ I. Mi i)

```

proof *safe*

fix i

assume h: limitin Product-metric.mtopology xn x sequentially

```

then show ∃ N. ∀ n ≥ N. (∀ i ∈ I. xn n i ∈ Mi i) ∧ (∀ i. i ∉ I → xn n i =
undefined)

```

by(simp only: Product-metric.limit-metric-sequentially) (metis PiE-E r(1))

```

then obtain N' where N': ∧ i n. i ∈ I ⇒ n ≥ N' ⇒ xn n i ∈ Mi i ∧ i n. i
∉ I ⇒ n ≥ N' ⇒ xn n i = undefined

```

by blast

assume i: i ∈ I

then interpret m: Metric-space Mi i di i

using Md-metric **by** blast

show limitin m.mtopology (λ n. xn n i) (x i) sequentially

unfolding m.limitin-metric eventually-sequentially

proof *safe*

show x i ∈ Mi i

using h i **by**(auto simp: Product-metric.limitin-metric)

next

fix ε :: real

assume 0 < ε

then obtain r ^ f i * ε > 0 **using** r **by** auto

then obtain N **where** N: ∧ n. n ≥ N ⇒ product-dist (xn n) x < r ^ f i * ε

using h(1) **by**(fastforce simp: Product-metric.limitin-metric eventually-sequentially)

show ∃ N. ∀ n ≥ N. xn n i ∈ Mi i ∧ di i (xn n i) (x i) < ε

proof(safe intro!: exI[where x=max N N'])

fix n

assume max N N' ≤ n

then have N ≤ n N' ≤ n

by auto

have di i (xn n i) (x i) ≤ (1 / r) ^ f i * product-dist (xn n) x

thm product-dist-geq[OF i gf-comp-id[OF i]]

```

using h i N'[OF - ⟨N' ≤ n⟩] by(auto intro!: product-dist-geq[OF i gf-comp-id[OF
i]] dest: Product-metric.limitin-mspace)

```

```

also have ... < (1 / r) ^ f i * r ^ f i * ε
  using N[OF <N ≤ n>] r by auto
also have ... ≤ ε
  by (simp add: <0 < ε> power-one-over)
finally show di i (xn n i) (x i) < ε .
qed(use N' h i in auto)
qed
next
fix N'
assume N': ∀ n ≥ N'. (∀ i ∈ I. xn n i ∈ Mi i) ∧ (∀ i. i ∉ I → xn n i = undefined)
assume h: ∀ i ∈ I. limitin (Metric-space.mtopology (Mi i) (di i)) (λ n. xn n i) (x i)
sequentially x ∈ (ΠE i ∈ I. Mi i)
show limitin Product-metric.mtopology xn x sequentially
  unfolding Product-metric.limitin-metric eventually-sequentially
proof safe
  fix ε
  assume he: (0 :: real) < ε
  then have 0 < ε * ((1 - r) / K) using r K-pos by auto
  hence ∃ k. r ^ k < ε * ((1 - r) / K)
    using r(2) real-arch-pow-inv by blast
  then obtain l where r ^ l < ε * ((1 - r) / K) by auto
  hence hk: r ^ l / (1 - r) * K < ε
    using mult-imp-div-pos-less[OF divide-pos-pos[OF - K-pos, of 1 - r]] r(2) by
simp
  hence hke: 0 < ε - r ^ l / (1 - r) * K by auto
  consider l = 0 | 0 < l by auto
  then show ∃ N. ∀ n ≥ N. xn n ∈ (ΠE i ∈ I. Mi i) ∧ product-dist (xn n) x < ε
proof cases
  case 1
  then have he2: 1 / (1 - r) * K < ε using hk by auto
  show ?thesis
    using order.strict-trans1[OF product-dist-leqr he2] N'
    by (auto intro!: exI[where x = N'])
next
  case 2
  with hke have 0 < 1 / real l * (ε - r ^ l / (1 - r) * K) by auto
  hence ∀ i ∈ I. ∃ N. ∀ n ≥ N. di i (xn n i) (x i) < 1 / real l * (ε - r ^ l / (1 - r) * K)
    using h by (meson Md-metric Metric-space.limit-metric-sequentially)
  then obtain N where hn:
  ∧ i n. i ∈ I ⇒ n ≥ N i ⇒ di i (xn n i) (x i) < 1 / real l * (ε - r ^ l / (1 - r) * K)
    by metis
  show ?thesis
proof (safe intro!: exI[where x = max (Sup {N (g n) | n. n < l}) N'])
  fix n
  assume max (⋂ {N (g n) | n. n < l}) N' ≤ n
  then have hsup: ⋂ {N (g n) | n. n < l} ≤ n and N'n: N' ≤ n
    by auto
  have product-dist (xn n) x = (∑ m. if g m ∈ I then r ^ m * di (g m) (xn n
(g m)) (x (g m)) else 0)

```

```

    using  $N' N'n h$  by(auto simp: product-dist-def)
    also have ... =  $(\sum m. \text{if } g(m+l) \in I \text{ then } r^{\wedge}(m+l)*di(g(m+l))$ 
 $(xn\ n(g(m+l))) (x(g(m+l))) \text{ else } 0) + (\sum m<l. \text{if } g\ m \in I \text{ then } r^{\wedge}m * di$ 
 $(g\ m) (xn\ n(g\ m)) (x(g\ m)) \text{ else } 0)$ 
    by(auto intro!: suminf-split-initial-segment)
    also have ...  $\leq r^{\wedge}l/(1-r)*K + (\sum m<l. \text{if } g\ m \in I \text{ then } r^{\wedge}m * di(g\ m)$ 
 $(xn\ n(g\ m)) (x(g\ m)) \text{ else } 0)$ 
    proof -
      have  $(\sum m. \text{if } g(m+l) \in I \text{ then } r^{\wedge}(m+l)*di(g(m+l)) (xn\ n(g$ 
 $(m+l))) (x(g(m+l))) \text{ else } 0) \leq (\sum m. r^{\wedge}(m+l)*K)$ 
      using di-bounded  $N' r K$ -pos by(auto intro!: suminf-le summable-ignore-initial-segment)
      also have ... =  $r^{\wedge}l/(1-r)*K$ 
      by(rule nsum-of-rK)
      finally show ?thesis by auto
    qed
    also have ...  $\leq r^{\wedge}l / (1 - r)*K + (\sum m<l. \text{if } g\ m \in I \text{ then } di(g\ m) (xn$ 
 $n(g\ m)) (x(g\ m)) \text{ else } 0)$ 
    proof -
      have  $(\sum m<l. \text{if } g\ m \in I \text{ then } r^{\wedge}m * di(g\ m) (xn\ n(g\ m)) (x(g\ m))$ 
 $\text{ else } 0) \leq (\sum m<l. \text{if } g\ m \in I \text{ then } di(g\ m) (xn\ n(g\ m)) (x(g\ m)) \text{ else } 0)$ 
      using di-bounded di-nonneg  $r$  by(auto intro!: sum-mono simp: mult-left-le-one-le
      power-le-one)
      thus ?thesis by simp
    qed
    also have ...  $< r^{\wedge}l / (1 - r)*K + (\sum m<l. 1 / \text{real } l * (\varepsilon - r^{\wedge}l/(1-r)*K))$ 
    proof -
      have  $(\sum m<l. \text{if } g\ m \in I \text{ then } di(g\ m) (xn\ n(g\ m)) (x(g\ m)) \text{ else } 0) <$ 
 $(\sum m<l. 1 / \text{real } l * (\varepsilon - r^{\wedge}l/(1-r)*K))$ 
      proof(rule sum-strict-mono-ex1)
        show  $\forall p \in \{..<l\}. (\text{if } g\ p \in I \text{ then } di(g\ p) (xn\ n(g\ p)) (x(g\ p)) \text{ else } 0)$ 
 $\leq 1 / \text{real } l * (\varepsilon - r^{\wedge}l / (1 - r)*K)$ 
        proof -
          have  $0 \leq (\varepsilon - r^{\wedge}l * K / (1 - r)) / \text{real } l$ 
          using hke by auto
          moreover {
            fix  $p$ 
            assume  $p < l$   $g\ p \in I$ 
            then have  $N(g\ p) \in \{N(g\ n) \mid n. n < l\}$ 
            by auto
            from le-cSup-finite[OF - this] hsup have  $N(g\ p) \leq n$ 
            by auto
            hence  $di(g\ p) (xn\ n(g\ p)) (x(g\ p)) \leq (\varepsilon - r^{\wedge}l * K / (1 - r)) /$ 
 $\text{real } l$ 
            using hn[OF  $\langle g\ p \in I \rangle$ , of  $n$ ] by simp
          }
        ultimately show ?thesis
        by auto
      qed
    next

```

```

      show  $\exists a \in \{..<l\}$ . (if  $g a \in I$  then  $di (g a) (xn n (g a)) (x (g a))$  else 0)
    < 1 / real l * ( $\varepsilon - r \wedge l / (1 - r) * K$ )
  proof -
    have  $0 < (\varepsilon - r \wedge l * K / (1 - r)) / \text{real } l$ 
      using hke 2 by auto
    moreover {
      assume  $g 0 \in I$ 
      have  $N (g 0) \in \{N (g n) \mid n. n < l\}$ 
        using 2 by auto
      from le-cSup-finite[OF - this] hsup have  $N (g 0) \leq n$ 
        by auto
      hence  $di (g 0) (xn n (g 0)) (x (g 0)) < (\varepsilon - r \wedge l * K / (1 - r)) /$ 
    }
    using hn[OF  $\langle g 0 \in I \rangle$ , of n] by simp
  }
  ultimately show ?thesis
    by(auto intro!: bexI[where x=0] simp: 2)
  qed
  qed simp
  thus ?thesis by simp
  qed
  also have ... =  $\varepsilon$ 
    using 2 by auto
  finally show product-dist (xn n) x <  $\varepsilon$  .
  qed(use N' in auto)
  qed
  qed (use N' h in auto)
qed(auto simp: Product-metric.limitin-metric)

lemma Product-metric-mtopology-eq: product-topology ( $\lambda i$ . Metric-space.mtopology
(Mi i) (di i)) I = Product-metric.mtopology
proof -
  have  $htospace: \bigwedge i. i \in I \implies \text{topspace } (\text{Metric-space.mtopology } (Mi i) (di i)) =$ 
Mi i
  by (simp add: Md-metric Metric-space.topspace-mtopology)
  hence  $htospace': (\prod_E i \in I. \text{topspace } (\text{Metric-space.mtopology } (Mi i) (di i))) =$ 
 $(\prod_E i \in I. Mi i)$  by auto
  consider  $I = \{\} \mid I \neq \{\}$  by auto
  then show ?thesis
  proof cases
    case 1
    interpret d: discrete-metric  $\{\lambda x. \text{undefined}\}$  .
    have product-dist =  $(\lambda x y. 0)$ 
      by standard+ (auto simp: product-dist-def 1)
    hence 2: Product-metric.mtopology = d.disc.mtopology
      by (metis 1 PiE-empty-domain Product-metric.open-in-mspace Product-metric.topspace-mtopology
d.mtopology-discrete-metric discrete-topology-unique singleton-iff)
    show ?thesis
      unfolding 2 by (simp add: product-topology-empty-discrete 1 d.mtopology-discrete-metric)
  end
end

```



```

next
  case I':2
  show ?thesis
  unfolding base-is-subbase[OF Product-metric.mtopology-base-in-balls,simplified
subbase-in-def] product-topology-def
  proof(rule topology-generated-by-eq)
    fix U
    assume U ∈ {Product-metric.mball a ε | a ε. a ∈ (ΠE i∈I. Mi i) ∧ 0 < ε}
    then obtain a ε where hu:
      U = Product-metric.mball a ε a ∈ (ΠE i∈I. Mi i) 0 < ε by auto
      have ∃ X. x ∈ (ΠE i∈I. X i) ∧ (ΠE i∈I. X i) ⊆ U ∧ (∀ i. openin
(Metric-space.mtopology (Mi i) (di i)) (X i)) ∧ finite {i. X i ≠ topspace (Metric-space.mtopology
(Mi i) (di i))} if x ∈ U for x
      proof -
        consider ε ≤ 1 / (1 - r) * K | 1 / (1 - r) * K < ε by fastforce
        then show ∃ X. x ∈ (ΠE i∈I. X i) ∧ (ΠE i∈I. X i) ⊆ U ∧ (∀ i.
openin (Metric-space.mtopology (Mi i) (di i)) (X i)) ∧ finite {i. X i ≠ topspace
(Metric-space.mtopology (Mi i) (di i))}
          proof cases
            case he2:1
              have 0 < (ε - product-dist a x)*((1-r)/ K) using r hu K-pos that hu
              by auto
              hence ∃ k. r~k < (ε - product-dist a x)*((1-r)/ K)
                using r(2) real-arch-pow-inv by blast
              then obtain k where r~k < (ε - product-dist a x)*((1-r)/ K) by auto
              hence hk:r~k / (1-r) * K < (ε - product-dist a x)
                using mult-imp-div-pos-less[OF divide-pos-pos[OF - K-pos,of 1-r]] r(2)
              by auto
              have hk': 0 < k apply(rule ccontr) using hk he2 Product-metric.nonneg[of
a x] by auto
              define ε' where ε' ≡ (1/(real k))*(ε - product-dist a x - r~k / (1-r) *
K)
              have hε': 0 < ε' using hk by(auto simp: ε'-def hk')
              define X where X ≡ (if finite I then (λi. if i ∈ I then Metric-space.mball
(Mi i) (di i) (x i) ε' else topspace (Metric-space.mtopology (Mi i) (di i))) else
(λi. if i ∈ I ∧ f i < k then Metric-space.mball (Mi i) (di i) (x i) ε' else topspace
(Metric-space.mtopology (Mi i) (di i))))
              show ?thesis
              proof(intro exI[where x=X] conjI)
                have x i ∈ Metric-space.mball (Mi i) (di i) (x i) ε' if i ∈ I for i
                using hu ⟨x ∈ U⟩ by (auto simp add: PiE-mem hε' Md-metric
Metric-spacecentre-in-mball-iff that)
                thus x ∈ (ΠE i∈I. X i)
                using hu that htopspace by(auto simp: X-def)
              next
              show (ΠE i∈I. X i) ⊆ U
              proof
                fix y
                assume y ∈ (ΠE i∈I. X i)

```

```

have  $\bigwedge i. X i \subseteq \text{topspace } (\text{Metric-space.mtopology } (Mi i) (di i))$ 
  by (simp add: Md-metric Metric-space.mball-subset-mspace X-def
htospace)
hence  $y \in (\prod_{E} i \in I. Mi i)$ 
  using htospace'  $\langle y \in (\prod_{E} i \in I. X i) \rangle$  by blast
have product-dist a  $y < \varepsilon$ 
proof -
  have product-dist a  $y \leq \text{product-dist a } x + \text{product-dist } x y$ 
    using Product-metric.triangle  $\langle y \in Pi_E I Mi \rangle hu(1)$  that by auto
  also have  $\dots < \text{product-dist a } x + (\varepsilon - \text{product-dist a } x)$ 
  proof -
    have product-dist  $x y < (\varepsilon - \text{product-dist a } x)$ 
    proof -
      have product-dist  $x y = (\sum n. \text{if } g n \in I \text{ then } r^{\wedge} n * di (g n) (x (g n)) (y (g n)) \text{ else } 0)$ 
        by (metis (no-types, lifting) hu(1) that  $\langle y \in (\prod_{E} i \in I. Mi i) \rangle$ 
Product-metric.in-mball product-dist-def suminf-cong)
      also have  $\dots = (\sum n. \text{if } g (n + k) \in I \text{ then } r^{\wedge} (n + k) * di (g (n + k)) (x (g (n + k))) (y (g (n + k))) \text{ else } 0) + (\sum n < k. \text{if } g n \in I \text{ then } r^{\wedge} n * di (g n) (x (g n)) (y (g n)) \text{ else } 0)$ 
        by (rule suminf-split-initial-segment) simp
      also have  $\dots \leq r^{\wedge} k / (1 - r) * K + (\sum n < k. \text{if } g n \in I \text{ then } r^{\wedge} n * di (g n) (x (g n)) (y (g n)) \text{ else } 0)$ 
        proof -
          have  $(\sum n. \text{if } g (n + k) \in I \text{ then } r^{\wedge} (n + k) * di (g (n + k)) (x (g (n + k))) (y (g (n + k))) \text{ else } 0) \leq (\sum n. r^{\wedge} (n + k) * K)$ 
            using di-bounded di-nonneg r K-pos by (auto intro!: suminf-le summable-ignore-initial-segment)
          also have  $\dots = r^{\wedge} k / (1 - r) * K$ 
            by (rule nsum-of-rK)
          finally show ?thesis by simp
        qed
      also have  $\dots < r^{\wedge} k / (1 - r) * K + (\varepsilon - \text{product-dist a } x - r^{\wedge} k / (1 - r) * K)$ 
        proof -
          have  $(\sum n < k. \text{if } g n \in I \text{ then } r^{\wedge} n * di (g n) (x (g n)) (y (g n)) \text{ else } 0) < (\sum n < k. \varepsilon')$ 
            proof (rule sum-strict-mono-ex1)
              show  $\forall l \in \{..<k\}. (\text{if } g l \in I \text{ then } r^{\wedge} l * di (g l) (x (g l)) (y (g l)) \text{ else } 0) \leq \varepsilon'$ 
                proof -
                  {
                    fix l
                    assume  $g l \in I l < k$ 
                    then interpret mbd: Metric-space Mi (g l) di (g l)
                      by (auto intro!: Md-metric)
                    have  $r^{\wedge} l * di (g l) (x (g l)) (y (g l)) \leq di (g l) (x (g l)) (y (g l))$ 
                      using r by (auto intro!: mult-right-mono [of  $r^{\wedge} l 1, OF -$ 

```

```

mbd.nonneg[of x (g l) y (g l),simplified] simp: power-le-one)
  also have ... < ε'
  proof -
    have y (g l) ∈ mbd.mball (x (g l)) ε'
    proof(cases finite I)
      case True
      then show ?thesis
        using PiE-mem[OF ⟨y ∈ (ΠE i∈I. X i)⟩ ⟨g l ∈ I⟩]
        by(simp add: X-def ⟨g l ∈ I⟩)
      next
      case False
      then show ?thesis
        using PiE-mem[OF ⟨y ∈ (ΠE i∈I. X i)⟩ ⟨g l ∈ I⟩]
gf-if-infinite(3)
        by(simp add: X-def ⟨g l ∈ I⟩ ⟨l < k⟩)
    qed
    thus ?thesis
      by auto
  qed
  finally have r ^ l * di (g l) (x (g l)) (y (g l)) ≤ ε' by simp
}
thus ?thesis
  by(auto simp: order.strict-implies-order[OF hε'])
qed
next
show ∃ a∈{..E i∈I. X i)⟩ ⟨g 0 ∈ I⟩]
      by(simp add: X-def ⟨g 0 ∈ I⟩)
    next
    case False
    then show ?thesis
      using PiE-mem[OF ⟨y ∈ (ΠE i∈I. X i)⟩ ⟨g 0 ∈ I⟩]
gf-if-infinite(3)
      by(simp add: X-def ⟨g 0 ∈ I⟩ ⟨0 < k⟩)
  qed
  hence r ^ 0 * di (g 0) (x (g 0)) (y (g 0)) < ε'
  by auto
}
thus (if g 0 ∈ I then r ^ 0 * di (g 0) (x (g 0)) (y (g 0)) else

```

$0) < \varepsilon'$

```

    using hε' by auto
    qed(use hk' in auto)
    qed simp
    also have ... = (ε - product-dist a x - r ^ k / (1 - r) * K)
    by(simp add: ε'-def hk')
    finally show ?thesis by simp
  qed
  finally show ?thesis by simp
  qed
  thus ?thesis by simp
  qed
  finally show ?thesis by auto
  qed
  thus y ∈ U
  by(simp add: hu(1) hu(2) ⟨y ∈ (ΠE i∈I. Mi i)⟩)
  qed
next
  have openin (Metric-space.mtopology (Mi i) (di i)) (Metric-space.mball
(Mi i) (di i) (x i) ε') if i ∈ I for i
  by (simp add: Md-metric Metric-space.openin-mball that)
  moreover have openin (Metric-space.mtopology (Mi i) (di i)) (topspace
(Metric-space.mtopology (Mi i) (di i))) for i
  by auto
  ultimately show ∀ i. openin (Metric-space.mtopology (Mi i) (di i)) (X
i)
  by(auto simp: X-def)
next
  show finite {i. X i ≠ topspace (Metric-space.mtopology (Mi i) (di i))}
  proof(cases finite I)
    case True
    then show ?thesis
    by(simp add: X-def)
  next
    case Iinf:False
    have finite {i ∈ I. f i < k}
    proof -
      have {i ∈ I. f i < k} = inv-into I f ' {..

```

```

      using *(1)[of p] by (auto simp: rev-image-eqI)
    qed
  next
    show inv-into I f ' {.. $k$ }  $\subseteq$  { $i \in I. f i < k$ }
      using *(2) bij-betw-inv-into[OF gf-if-infinite(1)[OF Iinf]]
      by (auto simp: bij-betw-def)
    qed
  qed
  also have finite ... by auto
  finally show ?thesis .
  qed
  thus ?thesis
    by(simp add: X-def Iinf)
  qed
  qed
next
  case 2
  then have  $U = (\prod_E i \in I. M i i)$ 
    unfolding hu(1) using order.strict-transI[OF product-dist-leqr, of  $\varepsilon$ ]
  hu(2)
    by auto
  also have ... =  $(\prod_E i \in I. \text{topspace } (Metric\text{-space.mtopology } (M i i) (d i i)))$ 
    using htopspace by auto
  finally have  $U = (\prod_E i \in I. \text{topspace } (Metric\text{-space.mtopology } (M i i) (d i i)))$  .
  thus ?thesis
    using that hu htopspace by(auto intro!: exI[where  $x = \lambda i. \text{topspace } (Metric\text{-space.mtopology } (M i i) (d i i))$ ])
  qed
  qed
  hence  $\exists X. \forall x \in U. x \in (\prod_E i \in I. X x i) \wedge (\prod_E i \in I. X x i) \subseteq U \wedge (\forall i. \text{openin } (Metric\text{-space.mtopology } (M i i) (d i i)) (X x i) \wedge \text{finite } \{i. X x i \neq \text{topspace } (Metric\text{-space.mtopology } (M i i) (d i i))\})$ 
    by(auto intro!: bchoice)
  then obtain X where  $\forall x \in U. x \in (\prod_E i \in I. X x i) \wedge (\prod_E i \in I. X x i) \subseteq U \wedge (\forall i. \text{openin } (Metric\text{-space.mtopology } (M i i) (d i i)) (X x i) \wedge \text{finite } \{i. X x i \neq \text{topspace } (Metric\text{-space.mtopology } (M i i) (d i i))\})$ 
    by auto
  hence hX:  $\bigwedge x. x \in U \implies x \in (\prod_E i \in I. X x i) \wedge x \in U \implies (\prod_E i \in I. X x i) \subseteq U \wedge x \in U \implies (\forall i. \text{openin } (Metric\text{-space.mtopology } (M i i) (d i i)) (X x i) \wedge x \in U \implies \text{finite } \{i. X x i \neq \text{topspace } (Metric\text{-space.mtopology } (M i i) (d i i))\})$ 
    by auto
  hence hXopen:  $\bigwedge x. x \in U \implies (\prod_E i \in I. X x i) \in \{\prod_E i \in I. X i i \mid X. (\forall i. \text{openin } (Metric\text{-space.mtopology } (M i i) (d i i)) (X i i) \wedge \text{finite } \{i. X i i \neq \text{topspace } (Metric\text{-space.mtopology } (M i i) (d i i))\})\}$ 
    by blast
  have  $U = (\bigcup \{(\prod_E i \in I. X x i) \mid x. x \in U\})$ 
    using hX(1,2) by blast

```

have *openin* (topology-generated-by $\{\Pi_E i \in I. X i \mid X. (\forall i. \text{openin} (\text{Metric-space.mtopology} (Mi i) (di i)) (X i)) \wedge \text{finite} \{i. X i \neq \text{topspace} (\text{Metric-space.mtopology} (Mi i) (di i))\}\}$) $\cup \{\{\Pi_E i \in I. X x i \mid x. x \in U\}\}$)
apply (rule *openin-Union*)
using *hXopen* **by** (auto *simp: openin-topology-generated-by-iff intro!: generate-topology-on.Basis*)
thus *openin* (topology-generated-by $\{\Pi_E i \in I. X i \mid X. (\forall i. \text{openin} (\text{Metric-space.mtopology} (Mi i) (di i)) (X i)) \wedge \text{finite} \{i. X i \neq \text{topspace} (\text{Metric-space.mtopology} (Mi i) (di i))\}\}$) U)
using $\langle U = (\cup \{\{\Pi_E i \in I. X x i \mid x. x \in U\}\}) \rangle$ **by** *simp*
next
fix U
assume $U \in \{\Pi_E i \in I. X i \mid X. (\forall i. \text{openin} (\text{Metric-space.mtopology} (Mi i) (di i)) (X i)) \wedge \text{finite} \{i. X i \neq \text{topspace} (\text{Metric-space.mtopology} (Mi i) (di i))\}\}$
then obtain X **where** *hX*:
 $U = (\Pi_E i \in I. X i) \wedge i. \text{openin} (\text{Metric-space.mtopology} (Mi i) (di i)) (X i)$
 $\text{finite} \{i. X i \neq \text{topspace} (\text{Metric-space.mtopology} (Mi i) (di i))\}$
by *auto*
have $\exists a \varepsilon. x \in \text{Product-metric.mball } a \varepsilon \wedge \text{Product-metric.mball } a \varepsilon \subseteq U$ **if**
 $x \in U$ **for** x
proof –
have $x \text{-intop}: x \in (\Pi_E i \in I. Mi i)$
unfolding *htopspace'* [*symmetric*] **using** *that hX(1) openin-subset[OF hX(2)]* **by** *auto*
define I' **where** $I' \equiv \{i. X i \neq \text{topspace} (\text{Metric-space.mtopology} (Mi i) (di i))\} \cap I$
then have $I': \text{finite } I' \ I' \subseteq I$ **using** *hX(3)* **by** *auto*
consider $I' = \{\} \mid I' \neq \{\}$ **by** *auto*
then show *?thesis*
proof *cases*
case 1
then have $\bigwedge i. i \in I \implies X i = \text{topspace} (\text{Metric-space.mtopology} (Mi i) (di i))$
by (auto *simp: I'-def*)
then have $U = (\Pi_E i \in I. Mi i)$
by (*simp add: PiE-eq hX(1) htopspace*)
thus *?thesis*
using 1 **that** **by** (auto *intro!: exI[where x=x] exI[where x=1]*)
next
case $I' \text{-nonempty}: 2$
hence $\bigwedge i. i \in I' \implies \text{openin} (\text{Metric-space.mtopology} (Mi i) (di i)) (X i)$
using *hX(2)* **by** (*simp add: I'-def*)
hence $\exists \varepsilon > 0. \text{Metric-space.mball} (Mi i) (di i) (x i) \varepsilon \subseteq (X i)$ **if** $i \in I'$ **for**
 i
using $I'(2)$ *Md-metric Metric-space.openin-mtopology PiE-mem* $\langle x \in U \rangle$
 $hX(1)$ *subset-eq* **that** **by** *blast*
then obtain $\varepsilon i'$ **where** $hei: \bigwedge i. i \in I' \implies \varepsilon i' i > 0 \ \bigwedge i. i \in I' \implies$
 $\text{Metric-space.mball} (Mi i) (di i) (x i) (\varepsilon i' i) \subseteq (X i)$
by *metis*

```

define  $\varepsilon$  where  $\varepsilon \equiv \text{Min } \{\varepsilon i' i \mid i. i \in I'\}$ 
have  $\varepsilon \text{min}$ :  $\bigwedge i. i \in I' \implies \varepsilon \leq \varepsilon i' i$ 
  using  $I'$  by(auto simp:  $\varepsilon$ -def intro!: Min.coboundedI)
have  $h\varepsilon$ :  $\varepsilon > 0$ 
  using  $I'$   $I'$ -nonempty Min-gr-iff[of  $\{\varepsilon i' i \mid i. i \in I'\}$  0] hei(1)
  by(auto simp:  $\varepsilon$ -def)
define  $n$  where  $n \equiv \text{Max } \{f i \mid i. i \in I'\}$ 
have  $\bigwedge i. i \in I' \implies f i \leq n$ 
  using  $I'$  by(auto intro!: Max.coboundedI[of  $\{f i \mid i. i \in I'\}$ ] simp: n-def)
hence  $hn2$ :  $\bigwedge i. i \in I' \implies (1 / r) \wedge f i \leq (1 / r) \wedge n$ 
  using  $r$  by auto
have  $h\varepsilon'$ :  $0 < \varepsilon * (r \wedge n)$  using  $h\varepsilon$   $r$  by auto
show ?thesis
proof(safe intro!: exI[where  $x=x$ ] exI[where  $x=\varepsilon*(r \wedge n)$ ])
  fix  $y$ 
  assume  $y \in \text{Product-metric.mball } x (\varepsilon * r \wedge n)$ 
  have  $y i \in X i$  if  $i \in I'$  for  $i$ 
  proof –
    interpret  $mi$ : Metric-space  $Mi i di i$ 
    using Md-metric that by(simp add: I'-def)
    have  $di i (x i) (y i) < \varepsilon i' i$ 
    proof –
      have  $di i (x i) (y i) \leq (1 / r) \wedge f i * \text{product-dist } x y$ 
      using that  $\langle y \in \text{Product-metric.mball } x (\varepsilon * r \wedge n) \rangle$  by(auto intro!:
product-dist-geq[of  $i, OF - gf-comp-id x-intop$ ] simp: I'-def)
      also have  $\dots \leq (1 / r) \wedge n * \text{product-dist } x y$ 
      by(rule mult-right-mono[OF hn2[OF that] Product-metric.nonneg])
      also have  $\dots < \varepsilon$ 
      using  $\langle y \in \text{Product-metric.mball } x (\varepsilon * r \wedge n) \rangle$   $r$ 
      by (simp add: pos-divide-less-eq power-one-over)
      also have  $\dots \leq \varepsilon i' i$ 
      by(rule  $\varepsilon \text{min}$ [OF that])
    finally show ?thesis .
  qed
hence  $(y i) \in mi.\text{mball } (x i) (\varepsilon i' i)$ 
  using  $\langle y \in \text{Product-metric.mball } x (\varepsilon * r \wedge n) \rangle$  x-intop that
  by(auto simp: I'-def)
thus ?thesis
  using hei[OF that] by auto
qed
moreover have  $y i \in X i$  if  $i \in I - I'$  for  $i$ 
  using that htopspace  $\langle y \in \text{Product-metric.mball } x (\varepsilon * r \wedge n) \rangle$ 
  by(auto simp: I'-def)
ultimately show  $y \in U$ 
  using  $\langle y \in \text{Product-metric.mball } x (\varepsilon * r \wedge n) \rangle$ 
  by(auto simp: hX(1))
qed(use x-intop h\varepsilon' in auto)
qed
qed

```

then obtain a where $\forall x \in U. \exists \varepsilon. x \in \text{Product-metric.mball } (a \ x) \ \varepsilon \wedge$
 $\text{Product-metric.mball } (a \ x) \ \varepsilon \subseteq U$
by *metis*
then obtain ε where $\text{hae}: \bigwedge x. x \in U \implies x \in \text{Product-metric.mball } (a \ x)$
 $(\varepsilon \ x) \wedge x. x \in U \implies \text{Product-metric.mball } (a \ x) \ (\varepsilon \ x) \subseteq U$
by *metis*
hence $\text{hae}' : \bigwedge x. x \in U \implies a \ x \in (\prod_{E \ i \in I. M \ i}) \wedge x. x \in U \implies 0 < \varepsilon \ x$
by *auto[1]* (*metis* *Product-metric.mball-eq-empty empty-iff hae(1) linorder-not-less*)
have $\text{openin } (\text{topology-generated-by } \{\text{Product-metric.mball } a \ \varepsilon \mid a \ \varepsilon. a \in (\prod_{E \ i \in I. M \ i}) \wedge 0 < \varepsilon\}) \cup \{\text{Product-metric.mball } (a \ x) \ (\varepsilon \ x) \mid x. x \in U\})$
using *Product-metric.openin-mball* $\langle \text{Product-metric.mtopology} = \text{topology-generated-by } \{\text{Product-metric.mball } a \ \varepsilon \mid a \ \varepsilon. a \in P_{i_E} \ I \ M \ i \wedge 0 < \varepsilon\} \rangle$ **by**
auto
moreover have $U = \cup \{\text{Product-metric.mball } (a \ x) \ (\varepsilon \ x) \mid x. x \in U\}$
using *hae* **by** (*auto simp del: Product-metric.in-mball*)
ultimately show $\text{openin } (\text{topology-generated-by } \{\text{Product-metric.mball } a \ \varepsilon \mid a \ \varepsilon. a \in (\prod_{E \ i \in I. M \ i}) \wedge 0 < \varepsilon\}) \ U$
by *simp*
qed
qed
qed

corollary *separable-Mi-separable-M:*

assumes $\bigwedge i. i \in I \implies \text{separable-space } (\text{Metric-space.mtopology } (M \ i) \ (d \ i))$
shows *separable-space* *Product-metric.mtopology*
by (*simp add: Product-metric.mtopology-eq[symmetric]*) *separable-countable-product*
assms I)

lemma *mcomplete-Mi-mcomplete-M:*

assumes $\bigwedge i. i \in I \implies \text{Metric-space.mcomplete } (M \ i) \ (d \ i)$
shows *Product-metric.mcomplete*

proof (*cases I = {}*)

case 1: *True*

interpret *d: discrete-metric* $\{\lambda x. \text{undefined}\}$.

have *2: product-dist =* $(\lambda x \ y. 0)$

by *standard+* (*auto simp: product-dist-def 1*)

show *?thesis*

apply (*simp add: Product-metric.mcomplete-def Product-metric.limitin-metric*
eventually-sequentially Product-metric.MCauchy-def)

apply (*simp add: 2*)

by (*auto simp: 1 intro!: exI[where x=($\lambda i. \text{undefined}$)]*)

next

assume *2: I \neq {}*

show *?thesis*

unfolding *Product-metric.mcomplete-def*

proof *safe*

fix *xn*

assume *h: Product-metric.MCauchy xn*

have **: Metric-space.MCauchy (M i) (d i) ($\lambda n. xn \ n \ i$)* **if** *hi: i \in I* **for** *i*


```

proof –
  interpret mi: Metric-space Mi i di i
    by(simp add: Md-metric hi)
  show mi.MCauchy ( $\lambda n. xn\ n\ i$ )
    unfolding mi.MCauchy-def
  proof safe
    show  $xn\ n\ i \in Mi\ i\ \text{for } n$ 
      using h hi by(auto simp: Product-metric.MCauchy-def)
  next
    fix  $\varepsilon$ 
    assume  $he:(0::real) < \varepsilon$ 
    then have  $0 < \varepsilon * r^{\wedge}(f\ i)$  using r by auto
    then obtain N where N:
       $\bigwedge n\ m. n \geq N \implies m \geq N \implies \text{product-dist } (xn\ n)\ (xn\ m) < \varepsilon * r^{\wedge}(f\ i)$ 
      using h Product-metric.MCauchy-def by fastforce
    show  $\exists N. \forall n\ n'. N \leq n \longrightarrow N \leq n' \longrightarrow di\ i\ (xn\ n\ i)\ (xn\ n'\ i) < \varepsilon$ 
    proof(safe intro!: exI[where x=N])
      fix n m
      assume  $n:n \geq N\ m \geq N$ 
      have  $di\ i\ (xn\ n\ i)\ (xn\ m\ i) \leq (1 / r)^{\wedge}(f\ i) * \text{product-dist } (xn\ n)\ (xn\ m)$ 
        by(rule product-dist-geq) (use h[simplified Product-metric.MCauchy-def])
      hi gf-comp-id[of i] in auto
      also have  $\dots < \varepsilon$ 
        using N[OF n] by (simp add: mult-imp-div-pos-less power-one-over r(1))
      finally show  $di\ i\ (xn\ n\ i)\ (xn\ m\ i) < \varepsilon$  .
    qed
  qed
  qed
  hence  $\forall i \in I. \exists x. \text{limitin } (Metric\text{-space.mtopology } (Mi\ i)\ (di\ i))\ (\lambda n. xn\ n\ i)\ x$ 
sequentially
    using Md-metric Metric-space.mcomplete-alt assms by blast
  then obtain x where  $hx:\bigwedge i. i \in I \implies \text{limitin } (Metric\text{-space.mtopology } (Mi\ i)\ (di\ i))\ (\lambda n. xn\ n\ i)\ (x\ i)$ 
sequentially
    by metis
  hence  $hx':(\lambda i \in I. x\ i) \in (\Pi_{E\ i \in I. Mi\ i})$ 
    by (simp add: Md-metric Metric-space.limit-metric-sequentially)
  thus  $\exists x. \text{limitin } Product\text{-metric.mtopology } xn\ x$ 
sequentially
    using h by(auto intro!: exI[where x=\lambda i \in I. x i] limitin-M-iff-limitin-Mi[THEN iffD2,of xn] simp: Product-metric.MCauchy-def hx) blast
  qed
qed

```

lemma *product-metricI*:

```

assumes  $0 < r < 1$  countable I  $\bigwedge i. i \in I \implies Metric\text{-space } (Mi\ i)\ (di\ i)$ 
  and  $\bigwedge i\ x\ y. 0 \leq di\ i\ x\ y \bigwedge i\ x\ y. di\ i\ x\ y \leq K\ 0 < K$ 
  shows Product-metric r I (to-nat-on I) (from-nat-into I) Mi di K
  using from-nat-into-to-nat-on-product-metric-pair[OF assms(3)] assms

```

by(simp add: Product-metric-def Metric-space-def)

lemma product-metric-natI:

assumes $0 < r$ $r < 1$ $\bigwedge n$. Metric-space (M n) (d n)
and $\bigwedge i x y$. $0 \leq d i x y$ $\bigwedge i x y$. $d i x y \leq K$ $0 < K$
shows Product-metric r UNIV id id M d K
using assms **by**(auto simp: Product-metric-def)

end

2 Abstract Polish Spaces

theory Abstract-Metrizable-Topology

imports Set-Based-Metric-Product

begin

2.1 Polish Spaces

definition Polish-space $X \equiv$ completely-metrizable-space $X \wedge$ separable-space X

lemma(in Metric-space) Polish-space-mtopology:

assumes mcomplete separable-space mtopology

shows Polish-space mtopology

by (simp add: assms completely-metrizable-space-mtopology Polish-space-def)

lemma

assumes Polish-space X

shows Polish-space-imp-completely-metrizable-space: completely-metrizable-space

X

and Polish-space-imp-metrizable-space: metrizable-space X

and Polish-space-imp-second-countable: second-countable X

and Polish-space-imp-separable-space: separable-space X

using assms **by**(auto simp: completely-metrizable-imp-metrizable-space Polish-space-def metrizable-space-separable-iff-second-countable)

lemma Polish-space-closedin:

assumes Polish-space X closedin $X A$

shows Polish-space (subtopology $X A$)

using assms **by**(auto simp: completely-metrizable-imp-metrizable-space Polish-space-def completely-metrizable-space-closedin second-countable-subtopology metrizable-space-separable-iff-second-countable)

lemma Polish-space-gdelta-in:

assumes Polish-space X gdelta-in $X A$

shows Polish-space (subtopology $X A$)

using assms **by**(auto simp: completely-metrizable-imp-metrizable-space Polish-space-def completely-metrizable-space-gdelta-in second-countable-subtopology metrizable-space-separable-iff-second-countable)

corollary Polish-space-openin:

assumes Polish-space X openin $X A$

shows *Polish-space (subtopology X A)*
by (*simp add: open-imp-gdelta-in assms Polish-space-gdelta-in*)

lemma *homeomorphic-Polish-space-aux:*
assumes *Polish-space X X homeomorphic-space Y*
shows *Polish-space Y*
using *assms by(simp add: homeomorphic-completely-metrizable-space-aux homeomorphic-separable-space Polish-space-def)*

corollary *homeomorphic-Polish-space:*
assumes *X homeomorphic-space Y*
shows *Polish-space X \longleftrightarrow Polish-space Y*
by (*meson assms homeomorphic-Polish-space-aux homeomorphic-space-sym*)

lemma *Polish-space-euclidean[simp]: Polish-space (euclidean :: ('a :: polish-space) topology)*
by (*simp add: completely-metrizable-space-euclidean Polish-space-def second-countable-imp-separable-space*)

lemma *Polish-space-countable[simp]:*
Polish-space (euclidean :: 'a :: {countable, discrete-topology} topology)
proof –
interpret *discrete-metric UNIV :: 'a set .*
have [*simp*]: *euclidean = disc.mtopology*
by (*metis discrete-topology-class.open-discrete discrete-topology-unique istopology-open mtopology-discrete-metric topology-inverse' topspace-euclidean*)
show *?thesis*
by (*auto simp: Polish-space-def intro!: disc.completely-metrizable-space-mtopology mcomplete-discrete-metric countable-space-separable-space*)
qed

lemma *Polish-space-discrete-topology: Polish-space (discrete-topology I) \longleftrightarrow countable I*
by (*simp add: completely-metrizable-space-discrete-topology Polish-space-def separable-space-discrete-topology*)

lemma *Polish-space-prod:*
assumes *Polish-space X and Polish-space Y*
shows *Polish-space (prod-topology X Y)*
using *assms by (simp add: completely-metrizable-space-prod-topology Polish-space-def separable-space-prod)*

lemma *Polish-space-product:*
assumes *countable I and $\bigwedge i. i \in I \implies$ Polish-space (S i)*
shows *Polish-space (product-topology S I)*
using *assms by(auto simp: separable-countable-product Polish-space-def completely-metrizable-space-product-topology)*

lemma(**in** *Product-metric*) *Polish-space I:*
assumes *$\bigwedge i. i \in I \implies$ separable-space (Metric-space.mtopology (M i) (d i))*

and $\bigwedge i. i \in I \implies \text{Metric-space.mcomplete } (Mi\ i) (di\ i)$
shows *Polish-space Product-metric.mtopology*
using *assms I* **by**(*auto simp: Polish-space-def Product-metric-mtopology-eq[symmetric]*
completely-metrizable-space-product-topology intro!: separable-countable-product Metric-space.completely-metrizable-space-mtopology Md-metric)

lemma(*in Sum-metric*) *Polish-spaceI*:
assumes *countable I*
and $\bigwedge i. i \in I \implies \text{separable-space } (\text{Metric-space.mtopology } (Mi\ i) (di\ i))$
and $\bigwedge i. i \in I \implies \text{Metric-space.mcomplete } (Mi\ i) (di\ i)$
shows *Polish-space Sum-metric.mtopology*
by(*auto simp: Polish-space-def intro!: separable-Mi-separable-M assms mcomplete-Mi-mcomplete-M Sum-metric.completely-metrizable-space-mtopology*)

lemma *compact-metrizable-imp-Polish-space*:
assumes *metrizable-space X compact-space X*
shows *Polish-space X*
proof –
obtain *d* **where** *Metric-space (topspace X) d Metric-space.mtopology (topspace X) d = X*
by (*metis assms(1) Metric-space.topspace-mtopology metrizable-space-def*)
thus *?thesis*
by (*metis Metric-space.compact-space-imp-separable assms(1) assms(2) compact-imp-locally-compact-space locally-compact-imp-completely-metrizable-space Polish-space-def*)
qed

2.2 Extended Reals and Non-Negative Extended Reals

lemma *Polish-space-ereal:Polish-space (euclidean :: ereal topology)*
proof(*rule homeomorphic-Polish-space-aux*)
show *Polish-space (subtopology euclideanreal $\{-1..1\}$)*
by (*auto intro!: Polish-space-closedin*)
next
define *f :: real \Rightarrow ereal*
where *f \equiv ($\lambda r. \text{if } r = -1 \text{ then } -\infty \text{ else if } r = 1 \text{ then } \infty \text{ else if } r \leq 0 \text{ then } \text{ereal } (1 - (1 / (1 + r))) \text{ else } \text{ereal } ((1 / (1 - r)) - 1)$)*
define *g :: ereal \Rightarrow real*
where *g \equiv ($\lambda r. \text{if } r \leq 0 \text{ then } \text{real-of-ereal } (\text{inverse } (1 - r)) - 1 \text{ else } 1 - \text{real-of-ereal } (\text{inverse } (1 + r))$)*
show *top-of-set $\{-1..1::\text{real}\}$ homeomorphic-space (euclidean :: ereal topology)*
unfolding *homeomorphic-space-def homeomorphic-maps-def continuous-map-iff-continuous*
proof(*safe intro!: exI[where x=f] exI[where x=g]*)
show *continuous-on $\{-1..1\}$ f*
unfolding *continuous-on-eq-continuous-within*
proof *safe*
fix *x :: real*
assume *x \in $\{-1..1\}$*
then consider *x = -1 | -1 < x x < 0 | x = 0 | 0 < x x < 1 | x = 1*

```

    by fastforce
  then show continuous (at x within {- 1..1}) f
  proof cases
    show  $-1 < x \implies x < 0 \implies ?thesis$ 
    by (simp add: at-within-Icc-at, intro isCont-cong[where f= $\lambda r. \text{ereal } (1 - (1 / (1 + r)))$ ]) and g=f, THEN iffD1, OF - continuous-on-interior[of  $\{-1 <..<0\}$ ])
    (auto simp: at-within-Icc-at eventually-nhds f-def intro!: exI[where x= $\{-1 <..<0\}$ ])
    continuous-on-divide continuous-on-ereal continuous-on-diff continuous-on-add
  next
    have *: isCont ( $\lambda r. \text{if } r \leq 0 \text{ then } \text{ereal } (1 - (1 / (1 + r))) \text{ else } \text{ereal } ((1 / (1 - r)) - 1)$ ) 0
    by (rule isCont-If-ge) (auto simp add: continuous-within intro!: continuous-on-Icc-at-leftD[where a= $-(1 / 2)$  and b=0 and f= $\lambda r::\text{real}. 1 - (1 / (1 + r))$ ,simplified] continuous-on-Icc-at-rightD[where a=0 and b= $1 / 2$  and f= $\lambda r::\text{real}. (1 / (1 - r)) - 1$ ,simplified] continuous-on-diff continuous-on-divide continuous-on-add)
    show ?thesis if x:x = 0
    unfolding x at-within-Icc-at[of  $-1 :: \text{real } 0 1$ ,simplified]
    by (rule isCont-cong[THEN iffD1, OF - *]) (auto simp: eventually-nhds f-def intro!: exI[where x= $\{-1 / 2 <..<1/2\}$ ])
  next
    show  $0 < x \implies x < 1 \implies ?thesis$ 
    by (simp add: at-within-Icc-at, intro isCont-cong[where f= $\lambda r. \text{ereal } ((1 / (1 - r)) - 1)$  and g=f, THEN iffD1, OF - continuous-on-interior[of  $\{0 <..<1\}$ ])
    (auto simp: at-within-Icc-at eventually-nhds f-def intro!: exI[where x= $\{0 <..<1\}$ ])
    continuous-on-divide continuous-on-ereal continuous-on-diff continuous-on-add
  next
    show ?thesis if x:x = -1
    unfolding x at-within-Icc-at-right[where a= $-1 :: \text{real}$  and b=1,simplified]
    continuous-within ereal-tendsto-simps(2)[symmetric]
    proof (subst tendsto-cong)
      show  $\forall_F r \text{ in } \text{at-right } (\text{ereal } (-1)). (f \circ \text{real-of-ereal}) r = 1 - (1 / (1 + r))$ 
      unfolding eventually-at-right[of  $\text{ereal } (-1) 0$ ,simplified]
      proof (safe intro!: exI[where x= $\text{ereal } (-1 / 2)$ ])
        fix y
        assume  $\text{ereal } (-1) < y < \text{ereal } (-1 / 2)$ 
        then obtain y' where y':  $y = \text{ereal } y' - 1 < y' < -1 / 2$ 
        by (metis ereal-real' less-ereal.simps(1) not-inftyI)
        show  $(f \circ \text{real-of-ereal}) y = 1 - 1 / (1 + y)$ 
        using y'(2,3) by (auto simp add: y'(1) f-def one-ereal-def)
      qed simp
    next
      have  $((\lambda r. 1 - 1 / (1 + r)) \longrightarrow -\infty) (\text{at-right } (\text{ereal } (-1)))$ 
      unfolding tendsto-MInfy
      proof safe
        fix r :: real
        consider  $r \geq 1 \mid r < 1$ 
        by argo

```

```

then show  $\forall_F x$  in at-right (ereal (- 1)).  $1 - 1 / (1 + x) < \text{ereal } r$ 
proof cases
  case [arith]:1
  show ?thesis
    unfolding eventually-at-right[of ereal (- 1) 0,simplified]
  proof(safe intro!: exI[where  $x=0$ ])
    fix  $y$ 
    assume ereal (- 1) <  $y$   $y < 0$ 
    then obtain  $y'$  where  $y':y = \text{ereal } y' - 1 < y'$   $y' < 0$ 
      using not-inftyI by force
    then have  $*:1 - 1 / (1 + y) < 1$ 
      by (simp add: one-ereal-def)
    show  $1 - 1 / (1 + y) < \text{ereal } r$ 
      by(rule order.strict-trans2[OF *]) (use 1 in auto)
    qed simp
  next
  case 2
  show ?thesis
    unfolding eventually-at-right[of ereal (- 1) 0,simplified]
  proof(safe intro!: exI[where  $x=\text{ereal } (1 / (1 - r) - 1)$ ])
    fix  $y$ 
    assume ereal (- 1) <  $y$   $y < \text{ereal } (1 / (1 - r) - 1)$ 
    then obtain  $y'$  where  $y':y = \text{ereal } y' - 1 < y'$   $y' < 1 / (1 - r) - 1$ 
      by (metis ereal-less-eq(3) ereal-real' linorder-not-le not-inftyI)
    have  $1 - 1 / (1 + y) = \text{ereal } (1 - 1 / (1 + y'))$ 
      by (metis ereal-divide ereal-minus(1) one-ereal-def order-less-irrefl
plus-ereal.simps(1) real-0-less-add-iff y'(1) y'(2))
    also have ... < ereal  $r$ 
    proof -
      have  $1 + y' < 1 / (1 - r)$ 
        using  $y'$  by simp
      hence  $1 - r < 1 / (1 + y')$ 
        using  $y'$  2 by (simp add: less-divide-eq mult commute)
      thus ?thesis
        by simp
    qed
    finally show  $1 - 1 / (1 + y) < \text{ereal } r$  .
    qed(use 2 in auto)
  qed
qed
thus (( $\lambda r. 1 - 1 / (1 + r)$ )  $\longrightarrow f (- 1)$ ) (at-right (ereal (- 1)))
  by(simp add: f-def)
qed
next
show ?thesis if  $x:x = 1$ 
  unfolding  $x$  at-within-Icc-at-left[where  $a=- 1 :: \text{real}$  and  $b=1$ ,simplified]
continuous-within ereal-tendsto-simps(1)[symmetric]
proof(subst tendsto-cong)
  show  $\forall_F r$  in at-left (ereal 1). ( $f \circ \text{real-of-ereal}$ )  $r = (1 / (1 - r)) - 1$ 

```

```

unfolding eventually-at-left[of 0 ereal 1,simplified]
proof(safe intro!: exI[where x=ereal (1 / 2)])
  fix y
  assume ereal (1 / 2) < y y < ereal 1
  then obtain y' where y':y = ereal y' 1 / 2 < y' y' < 1
    using ereal-less-ereal-Ex by force
  show (f ∘ real-of-ereal) y = 1 / (1 - y) - 1
    using y'(2,3) by(auto simp add: y'(1) f-def one-ereal-def)
qed simp
next
have ((λr. (1 / (1 - r)) - 1) → ∞) (at-left (ereal 1))
  unfolding tendsto-PInfty
proof safe
  fix r :: real
  consider r ≤ - 1 | - 1 < r
    by argo
  then show ∀F x in at-left (ereal 1). (1 / (1 - x)) - 1 > ereal r
proof cases
  case [arith]:1
  show ?thesis
    unfolding eventually-at-left[of 0 ereal 1,simplified]
proof(safe intro!: exI[where x=0])
  fix y
  assume 0 < y y < ereal 1
  then obtain y' where y':y = ereal y' 0 < y' y' < 1
    using not-inftyI by force
  then have *:(1 / (1 - y)) - 1 > ereal (- 1)
    by (simp add: one-ereal-def)
  show ereal r < 1 / (1 - y) - 1
    by(rule order.strict-trans1[OF - *]) (use 1 in auto)
qed simp
next
case 2
  show ?thesis
    unfolding eventually-at-left[of 0 ereal 1,simplified]
proof(safe intro!: exI[where x=ereal (1 - 1 / (1 + r))])
  fix y
  assume ereal (1 - 1 / (1 + r)) < y y < ereal 1
  then obtain y' where y':y = ereal y' 1 - 1 / (1 + r) < y' y' < 1
    by (metis ereal-less-eq(3) ereal-real' linorder-not-le not-inftyI)
  have ereal r < ereal (1 / (1 - y') - 1)
proof -
  have 1 - y' < 1 / (r + 1)
    using y'(2) by argo
  hence r + 1 < 1 / (1 - y')
    using y' 2 by (simp add: less-divide-eq mult commute)
  thus ?thesis
    by simp
qed

```

```

      also have ... = 1 / (1 - y) - 1
      by (metis diff-gt-0-iff-gt ereal-divide ereal-minus(1) less-numeral-extra(3))
one-ereal-def y'(1) y'(3))
      finally show ereal r < 1 / (1 - y) - 1 .
      qed(use 2 in auto)
    qed
  qed
  thus ((λr. 1 / (1 - r) - 1) ⟶ f 1) (at-left (ereal 1))
  by(simp add: f-def)
  qed
  qed
  qed
next
show continuous-map euclidean (top-of-set {- 1..1}) g
proof(safe intro!: continuous-map-into-subtopology)
  show continuous-map euclidean euclideanreal g
  unfolding Abstract-Topology.continuous-map-iff-continuous2 continuous-on-eq-continuous-within
  proof safe
    fix x :: ereal
    consider x = - ∞ | - ∞ < x x < 0 | x = 0 | 0 < x x < ∞ | x = ∞
    by fastforce
    then show isCont g x
    proof cases
      assume x: - ∞ < x x < 0
      then obtain x' where x': x = ereal x' x' < 0
      by (metis ereal-inf-ty-less(2) ereal-less-ereal-Ex zero-ereal-def)
      show ?thesis
      proof(subst isCont-cong)
        have [simp]: isCont ((-) 1) x
        proof -
          have *: isCont (λx. ereal (real-of-ereal 1 - real-of-ereal x)) x
          using x' by(auto simp add: continuous-at-iff-ereal[symmetric,simplified
comp-def] intro!: continuous-diff continuous-at-of-ereal)
          have **: ereal (x' - 1) < x ⟹ x < 0 ⟹ ereal (1 - real-of-ereal x)
= ereal 1 - x for x
          by (metis ereal-minus(1) less-ereal.simps(2) less-ereal.simps(3)
real-of-ereal.elims)
          show ?thesis
          apply(rule isCont-cong[THEN iffD1,OF - *])
          using x'(2) ** by(auto simp: eventually-nhds x'(1) one-ereal-def
intro!: exI[where x={x-1<..<0}])
        qed
        have *: abs (1 - x) ≠ ∞ 1 - x ≠ 0
        using x'(2) by(auto simp add: x'(1) one-ereal-def)
        show isCont (λr. real-of-ereal (inverse (1 - r)) - 1) x
        using x * by(auto intro!: continuous-diff continuous-divide isCont-o2[OF
- continuous-at-of-ereal])
      next
      show ∀F x in nhds x. g x = real-of-ereal (inverse (1 - x)) - 1

```



```

      using x(2) by(auto simp: eventually-nhds x'(1) g-def one-ereal-def
intro!: exI[where x={x-1<.. $0$ }]
    qed
  next
    assume x: $\infty > x > 0$ 
    then obtain x' where x': $x = \text{ereal } x' \ x' > 0$ 
      by (metis ereal-less(2) less-ereal.elims(2) less-ereal.simps(2))
    show ?thesis
    proof(subst isCont-cong)
      have [simp]: isCont ((+) 1) x
      proof -
        have *:isCont ( $\lambda x. \text{ereal } (\text{real-of-ereal } 1 + \text{real-of-ereal } x)$ ) x
        using x' by(auto simp add: continuous-at-iff-ereal[symmetric,simplified
comp-def] intro!: continuous-add continuous-at-of-ereal)
        have **:  $0 < x \implies x < \text{ereal } (x' + 1) \implies \text{ereal } (1 + \text{real-of-ereal } x)$ 
        =  $\text{ereal } 1 + x$  for x
          using ereal-less-ereal-Ex by auto
        show ?thesis
        apply(rule isCont-cong[THEN iffD1,OF - *])
          using x'(2) ** by(auto simp: eventually-nhds x'(1) one-ereal-def
intro!: exI[where x={ $0 < .. < x + 1$ }]
        qed
      have real-of-ereal (1 + x)  $\neq 0$ 
      using x' by auto
      thus isCont ( $\lambda r. 1 - \text{real-of-ereal } (\text{inverse } (1 + r))$ ) x
      using x by(auto intro!: continuous-diff continuous-divide isCont-o2[OF
- continuous-at-of-ereal])
    next
      show  $\forall_F x$  in nhds x.  $g \ x = 1 - \text{real-of-ereal } (\text{inverse } (1 + x))$ 
      using x(2) by(auto simp: eventually-nhds x'(1) g-def one-ereal-def
intro!: exI[where x={ $0 < .. < x + 1$ }]
    qed
  next
    show isCont g x if x: $x = - \infty$ 
      unfolding x
    proof(safe intro!: continuous-at-sequentiallyI)
      fix u :: nat  $\Rightarrow$  ereal
      assume u: $u \longrightarrow - \infty$ 
      show ( $\lambda n. g \ (u \ n)$ )  $\longrightarrow g \ (- \infty)$ 
        unfolding LIMSEQ-def
      proof safe
        fix r :: real
        assume r[arith]:  $r > 0$ 
        obtain no where no:  $\bigwedge n. n \geq no \implies u \ n < \text{ereal } (\text{min } (1 - 1 / r)$ 
0)
          using u unfolding tendsto-MInfty eventually-sequentially by blast
        show  $\exists no. \forall n \geq no. \text{dist } (g \ (u \ n)) \ (g \ (- \infty)) < r$ 
        proof(safe intro!: exI[where x=no])
          fix n

```

```

assume n:n ≥ no
have r0:1 - min (ereal (1 - 1 / r)) (ereal 0) > 0
  by (simp add: ereal-diff-gr0 min.strict-coboundedI2)
have u1:1 - u n > 0
  by (metis ereal-0-less-1 ereal-diff-gr0 ereal-min linorder-not-le
min.strict-coboundedI2 n no order-le-less-trans order-less-not-sym zero-ereal-def)
have real-of-ereal (inverse (1 - u n)) < r
proof -
  have real-of-ereal (inverse (1 - u n)) < real-of-ereal (inverse (1 -
ereal (min (1 - 1 / r) 0)))
  proof(safe intro!: ereal-less-real-iff[THEN iffD2])
  have ereal (real-of-ereal (inverse (1 - u n))) = inverse (1 - u n)
  by(rule ereal-real') (use no[OF n] u1 in auto)
  also have ... < inverse (1 - ereal (min (1 - 1 / r) 0))
  apply(rule ereal-inverse-antimono-strict)
  using no[OF n] apply(simp add: ereal-diff-positive min.coboundedI2)
  by (metis (no-types, lifting) no[OF n] ereal-add-uminus-conv-diff
ereal-eq-minus-iff ereal-less-minus-iff ereal-minus-less-minus ereal-times(1) ereal-times(3))
  finally show ereal (real-of-ereal (inverse (1 - u n))) < inverse
(1 - ereal (min (1 - 1 / r) 0)) .
  qed(use r0 in auto)
  also have ... ≤ r
  by(cases r ≥ 1) (auto simp add: real-of-ereal-minus)
  finally show real-of-ereal (inverse (1 - u n)) < r .
qed
thus dist (g (u n)) (g (- ∞)) < r
  using u1 no[OF n] by(auto simp: g-def zero-ereal-def dist-real-def)
qed
qed
qed
next
show isCont g x if x:x = ∞
  unfolding x
proof(safe intro!: continuous-at-sequentiallyI)
  fix u :: nat ⇒ ereal
  assume u:u → ∞
  show (λn. g (u n)) → g ∞
  unfolding LIMSEQ-def
proof safe
  fix r :: real
  assume r[arith]: r > 0
  obtain no where no: ∧n. n ≥ no ⇒ u n > ereal (max (1 / r - 1)
0)

  using u unfolding tendsto-PInfty eventually-sequentially by blast
show ∃no. ∀n≥no. dist (g (u n)) (g ∞) < r
proof(safe intro!: exI[where x=no])
  fix n
  assume n:n ≥ no
  have u0: 1 + u n > 0

```

```

      using no[OF n] by simp (metis add-nonneg-pos zero-ereal-def
zero-less-one-ereal)
      have  $| - \text{real-of-ereal} (\text{inverse} (1 + u n)) | < r$ 
      proof -
        have  $| - \text{real-of-ereal} (\text{inverse} (1 + u n)) | < | - (\text{inverse} (1 + \max
(1 / r - 1) 0)) |$ 
          unfolding abs-real-of-ereal abs-minus
          proof (safe intro!: real-less-ereal-iff[THEN iffD2])
            have  $|\text{inverse} (1 + u n)| < \text{inverse} (1 + \text{ereal} (\max (1 / r - 1)
0))$ 
              using no[OF n] u0 by (simp add: ereal-add-strict-mono
ereal-inverse-antimono-strict inverse-ereal-ge0I le-max-iff-disj order-less-imp-le u0)
              also have  $\dots = \text{ereal} |\text{inverse} (1 + \max (1 / r - 1) 0)|$ 
              by (auto simp: abs-ereal.simps(1)[symmetric] ereal-max[symmetric]
simp del: abs-ereal.simps(1) ereal-max)
              finally show  $|\text{inverse} (1 + u n)| < \text{ereal} |\text{inverse} (1 + \max (1 /
r - 1) 0)| .$ 
            qed auto
            also have  $\dots = \text{inverse} (1 + \max (1 / r - 1) 0)$ 
            by auto
            also have  $\dots \leq r$ 
            by (cases  $r \leq 1$ ) auto
            finally show ?thesis .
          qed
        thus  $\text{dist} (g (u n)) (g \infty) < r$ 
        using no[OF n] by (auto simp: g-def dist-real-def zero-ereal-def)
      qed
    qed
  qed
next
show isCont g x if  $x : x = 0$ 
  unfolding x g-def
  proof (safe intro!: isCont-If-ge)
    have  $((\lambda x. \text{real-of-ereal} (1 - x)) \longrightarrow 1) (\text{at-left } 0)$ 
    proof (subst tendsto-cong)
      show  $((\lambda x. 1 - \text{real-of-ereal } x) \longrightarrow 1) (\text{at-left } 0)$ 
      by (auto intro!: tendsto-diff[where a=1 and b=0,simplified] simp:
zero-ereal-def)
    next
      show  $\forall_F x \text{ in } \text{at-left } 0. \text{real-of-ereal} (1 - x) = 1 - \text{real-of-ereal } x$ 
      by (auto simp: eventually-at-left[where y=- 1 and x=0::ereal,simplified]
real-of-ereal-minus ereal-uminus-eq-reorder intro!: exI[where x=-1])
    qed
    thus continuous (at-left 0)  $(\lambda x. \text{real-of-ereal} (\text{inverse} (1 - x)) - 1)$ 
    unfolding continuous-within
    by (auto intro!: tendsto-diff[where a = 1 and b=1,simplified]
tendsto-divide[where a=1 and b=1,simplified])
  next
    have  $((\lambda x. \text{real-of-ereal} (1 + x)) \longrightarrow 1) (\text{at-right } 0)$ 

```

```

proof(subst tendsto-cong)
  show (( $\lambda x. 1 + \text{real-of-ereal } x \longrightarrow 1$ ) (at-right 0))
    by(auto intro!: tendsto-add[where a=1 and b=0,simplified] simp:
zero-ereal-def)
  next
    show  $\forall_F x$  in at-right 0.  $\text{real-of-ereal } (1 + x) = 1 + \text{real-of-ereal } x$ 
    by(auto simp: eventually-at-right[where y=1 and x=0::ereal,simplified]
real-of-ereal-add ereal-uminus-eq-reorder intro!: exI[where x=1])
    qed
    thus (( $\lambda x. 1 - \text{real-of-ereal } (\text{inverse } (1 + x)) \longrightarrow \text{real-of-ereal } (\text{inverse } (1 - 0)) - 1$ ) (at-right 0))
    by (auto intro!: tendsto-diff[where a = 1 and b=1,simplified]
tendsto-divide[where a=1 and b=1,simplified])
    qed
  qed
qed
next
fix x :: ereal
consider x = -  $\infty$  | -  $\infty$  < x x  $\leq$  0 | 0 < x x <  $\infty$  | x =  $\infty$ 
  by fastforce
then show g x  $\in$  {- 1..1}
proof cases
  assume -  $\infty$  < x x  $\leq$  0
  then obtain x' where x = ereal x' x'  $\leq$  0
    by (metis dual-order.refl ereal-less-ereal-Ex order-less-le zero-ereal-def)
  then show ?thesis
  by(auto simp: g-def real-of-ereal-minus intro!: pos-divide-le-eq[THEN iffD2])
next
  assume 0 < x x <  $\infty$ 
  then obtain x' where x = ereal x' x' > 0
    by (metis ereal-less(2) less-ereal.elims(2) order-less-le)
  then show ?thesis
  by(auto simp: g-def real-of-ereal-add inverse-eq-divide intro!: pos-divide-le-eq[THEN
iffD2])
  qed(auto simp: g-def)
qed
next
fix x :: ereal
consider x = -  $\infty$  | -  $\infty$  < x x  $\leq$  0 | 0 < x x <  $\infty$  | x =  $\infty$ 
  by fastforce
then show f (g x) = x
proof cases
  assume -  $\infty$  < x x  $\leq$  0
  then obtain x' where x':x = ereal x' x'  $\leq$  0
    by (metis dual-order.refl ereal-less-ereal-Ex order-less-le zero-ereal-def)
  then have [arith]: 1 / (1 - x') - 1  $\leq$  0
    by simp
  show ?thesis
    using x' by(auto simp: g-def real-of-ereal-minus f-def)

```

```

next
  assume  $0 < x < \infty$ 
  then obtain  $x'$  where  $x':x = \text{ereal } x' \ x' > 0$ 
    by (metis ereal-less(2) less-ereal.elims(2) order-less-le)
  hence [arith]:  $1 - 1 / (x' + 1) \geq 0$ 
    by simp
  show ?thesis
    using  $x'$  by (simp add: g-def inverse-eq-divide f-def)
qed(auto simp: f-def g-def)
next
fix  $x :: \text{real}$ 
assume  $x \in \text{topspace } (\text{top-of-set } \{-1..1\})$ 
then consider  $x = -1 \mid -1 < x \leq 0 \mid 0 < x < 1 \mid x = 1$ 
  by fastforce
then show  $g(f x) = x$ 
  by cases (auto simp: f-def g-def real-of-ereal-minus real-of-ereal-add)
qed
qed

```

corollary *Polish-space-ennreal:Polish-space (euclidean :: ennreal topology)*

proof(rule *homeomorphic-Polish-space-aux*)

show *Polish-space (top-of-set {0::ereal..})*

using *Polish-space-closedin Polish-space-ereal* by fastforce

next

show *top-of-set {0::ereal..} homeomorphic-space (euclidean :: ennreal topology)*

by (auto intro!: $\text{exI}[\text{where } x = \text{e2ennreal}] \text{exI}[\text{where } x = \text{enn2ereal}] \text{simp: homeomorphic-space-def homeomorphic-maps-def enn2ereal-e2ennreal continuous-on-e2ennreal continuous-map-in-subtopology continuous-on-enn2ereal image-subset-iff}$)

qed

2.3 Continuous Embddings

abbreviation *Hilbert-cube-topology :: (nat \Rightarrow real) topology where*

Hilbert-cube-topology \equiv (product-topology ($\lambda n. \text{top-of-set } \{0..1\}$) UNIV)

lemma *topspace-Hilbert-cube: topspace Hilbert-cube-topology = ($\Pi_E x \in \text{UNIV}. \{0..1\}$)*

by simp

lemma *Polish-space-Hilbert-cube: Polish-space Hilbert-cube-topology*

by (auto intro!: *Polish-space-closedin Polish-space-product*)

abbreviation *Cantor-space-topology :: (nat \Rightarrow real) topology where*

Cantor-space-topology \equiv (product-topology ($\lambda n. \text{top-of-set } \{0,1\}$) UNIV)

lemma *topspace-Cantor-space:*

topspace Cantor-space-topology = ($\Pi_E x \in \text{UNIV}. \{0,1\}$)

by simp

lemma *Polish-space-Cantor-space: Polish-space Cantor-space-topology*

by(*auto intro!*: Polish-space-closedin Polish-space-product)

corollary *completely-metrizable-space-homeo-image-gdelta-in*:

assumes *completely-metrizable-space X completely-metrizable-space Y B* \subseteq *topspace Y X homeomorphic-space subtopology Y B*

shows *gdelta-in Y B*

using *assms completely-metrizable-space-eq-gdelta-in homeomorphic-completely-metrizable-space*
by *blast*

2.3.1 Embedding into Hilbert Cube

lemma *embedding-into-Hilbert-cube*:

assumes *metrizable-space X separable-space X*

shows $\exists A \subseteq$ *topspace Hilbert-cube-topology. X homeomorphic-space (subtopology Hilbert-cube-topology A)*

proof –

consider $X =$ *trivial-topology* | *topspace X* \neq $\{\}$ **by** *auto*

then show *?thesis*

proof *cases*

case *1*

then show *?thesis*

by(*auto intro!*: *exI*[**where** $x = \{\}$]] *simp: homeomorphic-empty-space-eq*)

next

case *S-ne:2*

then obtain U **where** U :*countable U dense-in X U U* \neq $\{\}$

using *assms(2)* **by**(*auto simp: separable-space-def2 dense-in-nonempty*)

obtain xn **where** xn : $\wedge n::nat. xn n \in U U =$ *range xn*

by (*metis U(1) U(3) from-nat-into range-from-nat-into*)

then have $xns: xn n \in$ *topspace X* **for** n

using *dense-in-subset[OF U(2)]* **by** *auto*

obtain d' **where** d' :*Metric-space (topspace X) d' Metric-space.mtopology (topspace X) d' = X*

by (*metis Metric-space.topspace-mtopology assms(1) metrizable-space-def*)

interpret ms' : *Metric-space topspace X d'* **by** *fact*

define d **where** $d = ms'.capped-dist (1/2)$

have d : *Metric-space.mtopology (topspace X) d = X* $\wedge x y. d x y < 1$

by(*simp add: d-def ms'.mtopology-capped-metric d'*) (*simp add: d-def ms'.capped-dist-def*)

interpret ms : *Metric-space topspace X d*

by (*simp add: d-def ms'.capped-dist*)

define f **where** $f \equiv (\lambda x n. d x (xn n))$

have f -*inj:inj-on f (topspace X)*

proof

fix $x y$

assume $xy: x \in$ *topspace X y \in topspace X f x = f y*

then have $\wedge n. d x (xn n) = d y (xn n)$ **by**(*auto simp: f-def dest: fun-cong*)

hence $d2: d x y \leq 2 * d x (xn n)$ **for** n

using *ms.triangle[OF xy(1) - xy(2), of xn n, simplified ms commute*[*of xn n y]] dense-in-subset[OF U(2)] xn(1)[of n]*

by *auto*

```

have d x y < ε if ε > 0 for ε
proof -
  have 0 < ε / 2 using that by simp
  then obtain n where d x (xn n) < ε / 2
    using ms.mdense-def2[of U,simplified d(1)] U(2) xy(1) xn(2) by blast
  with d2[of n] show ?thesis by simp
qed
hence d x y = 0
  by (metis ms.nonneg[of x y] dual-order.irreft order-neq-le-trans)
thus x = y
  using xy by simp
qed
have f-img: f ' topspace X ⊆ topspace Hilbert-cube-topology
  using d(2) ms.nonneg by(auto simp: topspace-Hilbert-cube f-def less-le-not-le)
have f-cont: continuous-map X Hilbert-cube-topology f
unfolding continuous-map-componentwise-UNIV f-def continuous-map-in-subtopology
proof safe
  show continuous-map X euclideanreal (λx. d x (xn k)) for k
  proof(rule continuous-map-eq[of - - mdist-set ms.Self {xn k}])
    show continuous-map X euclideanreal (mdist-set ms.Self {xn k})
    by (metis d(1) mdist-set-uniformly-continuous ms.mdist-Self ms.mspace-Self
mtopology-of-def mtopology-of-euclidean uniformly-continuous-imp-continuous-map)
  next
    fix x
    assume x ∈ topspace X
    then show mdist-set ms.Self {xn k} x = d x (xn k)
      by(auto simp: ms.mdist-set-Self xns)
  qed next
  show d x (xn k) ∈ {0..1} for x k
    using d(2) ms.nonneg by(auto simp: less-le-not-le)
  qed
hence f-cont': continuous-map X (subtopology Hilbert-cube-topology (f ' topspace
X)) f
  using continuous-map-into-subtopology by blast
obtain g where g: g ' (f ' topspace X) = topspace X ∧ x. x ∈ topspace X ⇒
g (f x) = x ∧ x. x ∈ f ' topspace X ⇒ f (g x) = x
  by (meson f-inj f-the-inv-into-f the-inv-into-f-eq the-inv-into-onto)
have g-cont: continuous-map (subtopology Hilbert-cube-topology (f ' topspace
X)) X g
proof -
  interpret m01: Submetric UNIV dist {0..1::real}
  by(simp add: Submetric-def Submetric-axioms-def Met-TC.Metric-space-axioms)
  have m01-eq: m01.sub.mtopology = top-of-set {0..1}
    using m01.mtopology-submetric by auto
  have m01-Polish: Polish-space m01.sub.mtopology
    by(auto simp: m01-eq intro!: Polish-space-closedin)
  interpret m01': Metric-space {0..1::real} λx y. if 0 ≤ x ∧ x ≤ 1 ∧ 0 ≤ y
  ∧ y ≤ 1 then dist x y else 0
    by(auto intro!: Metric-space-eq[OF m01.sub.Metric-space-axioms]) metric

```

```

have m01'-eq: m01'.mtopology = top-of-set {0..1}
by(auto intro!: Metric-space-eq-mtopology[OF m01.sub.Metric-space-axioms,simplified
m01-eq,symmetric]) metric
have dist x y ≤ 1 dist x y ≥ 0 if x ≥ 0 x ≤ 1 y ≥ 0 y ≤ 1 for x y :: real
using dist-real-def that by auto
then interpret ppm: Product-metric 1/2 UNIV :: nat set id id λ-. {0..1::real}
λ- x y. if 0 ≤ x ∧ x ≤ 1 ∧ 0 ≤ y ∧ y ≤ 1 then dist x y else 0 1
by(auto intro!: product-metric-natI Metric-space-eq[OF m01.sub.Metric-space-axioms]
simp: m01.sub commute)

have Hilbert-cube-eq: ppm.Product-metric.mtopology = Hilbert-cube-topology
by(simp add: ppm.Product-metric-mtopology-eq[symmetric] m01'-eq)
interpret f-S: Submetric ΠE x∈UNIV. {0..1} ppm.product-dist f ' topspace
X
by(auto simp: Submetric-def ppm.Product-metric.Metric-space-axioms Sub-
metric-axioms-def f-def order.strict-implies-order[OF d(2)])
have 1:subtopology Hilbert-cube-topology (f ' topspace X) = f-S.sub.mtopology
using Hilbert-cube-eq f-S.mtopology-submetric by auto
have continuous-map f-S.sub.mtopology ms.mtopology g
unfolding continuous-map-iff-limit-seq[OF f-S.sub.first-countable-mtopology]
proof safe
fix yn y
assume h: limitin f-S.sub.mtopology yn y sequentially
have h':limitin ppm.Product-metric.mtopology yn y sequentially
using f-S.limitin-submetric-iff h by blast
hence m01-conv:∧n. limitin m01'.mtopology (λi. yn i n) (y n) sequentially
y ∈ UNIV →E {0..1}
by(auto simp: ppm.limitin-M-iff-limitin-Mi)
have ∃N. ∀n≥N. ∃zn. yn n = f zn ∧ zn ∈ topspace X
using h g by(simp only: f-S.sub.limit-metric-sequentially) (meson imageE
ppm.K-pos)
then obtain N' zn where zn:∧n. n ≥ N' ⇒ f (zn n) = yn n ∧n. n ≥
N' ⇒ zn n ∈ topspace X
by metis
obtain z where z:f z = y z ∈ topspace X
using h f-S.sub.limitin-mspace by blast
show limitin ms.mtopology (λn. g (yn n)) (g y) sequentially
unfolding ms.limit-metric-sequentially
proof safe
fix ε :: real
assume he: 0 < ε
then have 0 < ε / 3 by simp
then obtain m where m:d z (xn m) < ε / 3
using ms.mdense-def2[of U,simplified d(1)] U(2) z(2) xn(2) by blast
have ∧e. e>0 ⇒ ∃N. ∀n≥N. yn n m ∈ {0..1} ∧ dist (yn n m) (y m)
< e
using m01-conv(1)[of m,simplified m01'.limit-metric-sequentially]
by fastforce
from this[OF ‹0 < ε / 3›] obtain N where ∧n. n ≥ N ⇒ |yn n m -

```



```

y m| < ε / 3 ∧ n. n ≥ N ⇒ yn n m ∈ {0..1}
  by(auto simp: dist-real-def)
  hence N: ∧ n. n ≥ N ⇒ yn n m < ε / 3 + y m
    by (metis abs-diff-less-iff add commute)
  have ∃ N. ∀ n ≥ N. zn n ∈ topspace X ∧ d (zn n) z < ε
  proof (safe intro!: exI [where x = max N N'])
    fix n
    assume max N N' ≤ n
    then have N ≤ n N' ≤ n
      by auto
    then have d (zn n) z ≤ f (zn n) m + d z (xn m)
    using ms.triangle[OF zn(2)[of n] xns[of m] z(2), simplified ms commute [of
xn m z]]
      by (auto simp: f-def)
    also have ... < ε / 3 + y m + d z (xn m)
      using N [OF ‹N ≤ n›] zn(1)[of n] ‹N' ≤ n› by simp
    also have ... = ε / 3 + d z (xn m) + d z (xn m)
      by (simp add: z(1)[symmetric] f-def)
    also have ... < ε
      using m by auto
    finally show d (zn n) z < ε .
  qed (use zn in auto)
  thus ∃ N. ∀ n ≥ N. g (yn n) ∈ topspace X ∧ d (g (yn n)) (g y) < ε
    by (metis dual-order.trans nle-le zn(1) z(1) g(2) [OF z(2)] g(2) [OF
zn(2)])
  qed (use g z in auto)
  qed
  hence continuous-map f-S.sub.mtopology ms.mtopology g
    by (auto simp: mtopology-of-def)
  thus ?thesis
    by (simp add: d(1) 1)
  qed
  show ?thesis
    using f-img g(2,3) f-cont' g-cont
  by (auto intro!: exI [where x = f ' topspace X] homeomorphic-maps-imp-homeomorphic-space [where
f=f and g=g] simp: homeomorphic-maps-def)
  qed
  qed

```

corollary embedding-into-Hilbert-cube-gdelta-in:
assumes Polish-space X
shows ∃ A. gdelta-in Hilbert-cube-topology A ∧ X homeomorphic-space (subtopology Hilbert-cube-topology A)
proof –
obtain A **where** h: A ⊆ topspace Hilbert-cube-topology X homeomorphic-space subtopology Hilbert-cube-topology A
using embedding-into-Hilbert-cube Polish-space-imp-metrizable-space Polish-space-imp-separable-space **assms** **by** blast
with completely-metrizable-space-homeo-image-gdelta-in [OF Polish-space-imp-completely-metrizable-space [O

assms] Polish-space-imp-completely-metrizable-space[OF Polish-space-Hilbert-cube]
h(1,2)]
show ?thesis
by blast
qed

2.3.2 Embedding from Cantor Space

lemma embedding-from-Cantor-space:

assumes Polish-space X uncountable (topspace X)
shows $\exists A. \text{gdelta-in } X A \wedge \text{Cantor-space-topology homeomorphic-space (subtopology } X A)$

proof –

obtain $U P$ **where** up : countable U openin $X U$ perfect-set $X P U \cup P = \text{topspace } X U \cap P = \{\}$ $\wedge a. a \neq \{\} \implies \text{openin (subtopology } X P) a \implies \text{uncountable } a$

using Cantor-Bendixon[OF Polish-space-imp-second-countable[OF *assms*(1)]]

by auto

have P : closedin $X P P \subseteq \text{topspace } X$ uncountable P

using countable-Un-iff[of $U P$] $up(1)$ *assms* $up(4)$

by(simp-all add: perfect-setD[OF $up(3)$])

then have pp : Polish-space (subtopology $X P$)

by (simp add: *assms*(1) Polish-space-closedin)

have $Ptop$: topspace (subtopology $X P) = P$

using $P(2)$ **by** auto

obtain U **where** U : countable U dense-in (subtopology $X P$) U

using Polish-space-def pp separable-space-def2 **by** auto

with uncountable-infinite[OF $P(3)$] $P(2)$

have infinite U

by (metis Metric-space.t1-space-mtopology $Ptop$ *assms*(1) completely-metrizable-space-def dense-in-infinite Polish-space-def t1-space-subtopology)

obtain d **where** Metric-space $P d$ **and** d : Metric-space.mtopology $P d = \text{subtopology } X P$ **and** mdc : Metric-space.mcomplete $P d$

by (metis Metric-space.topspace-mtopology $Ptop$ completely-metrizable-space-def Polish-space-def pp)

interpret md : Metric-space $P d$ **by** fact

define xn **where** $xn \equiv \text{from-nat-into } U$

have xn : bij-betw $xn UNIV U \wedge n m. n \neq m \implies xn n \neq xn m \wedge n. xn n \in U \wedge n. xn n \in P \text{md.mdense (range } xn)$

using bij-betw-from-nat-into[OF $U(1)$ $\langle \text{infinite } U \rangle$] dense-in-subset[OF $U(2)$] $d U(2)$ range-from-nat-into[OF infinite-imp-nonempty[OF $\langle \text{infinite } U \rangle$] $U(1)$]

by(auto simp add: xn -def $U(1)$ $\langle \text{infinite } U \rangle$ from-nat-into[OF infinite-imp-nonempty[OF $\langle \text{infinite } U \rangle$]])

have perfect:perfect-space md .mtopology

using d perfect-set-subtopology $up(3)$ **by** simp

define jn **where** $jn \equiv (\lambda n. \text{LEAST } i. i > n \wedge \text{md.mcball } (xn i) ((1/2) \hat{\sim} i) \subseteq \text{md.mball } (xn n) ((1/2) \hat{\sim} n) - \text{md.mball } (xn n) ((1/2) \hat{\sim} i))$

define kn **where** $kn \equiv (\lambda n. \text{LEAST } k. k > jn n \wedge \text{md.mcball } (xn k) ((1/2) \hat{\sim} k) \subseteq \text{md.mball } (xn n) ((1/2) \hat{\sim} jn n))$

have $d \times m : xn. \forall n n'. \exists m. m > n \wedge m > n' \wedge (1/2) \hat{\sim} (m-1) < d (xn n) (xn m)$

$\wedge d(xn\ n) (xn\ m) < (1/2) \wedge (Suc\ n')$
proof safe
fix $n\ n'$
have $hinfin':infinite (md.mball\ x\ e \cap (range\ xn))$ **if** $x \in P\ e > 0$ **for** $x\ e$
proof
assume $h-fin:finite (md.mball\ x\ e \cap range\ xn)$
have $h-nen:md.mball\ x\ e \cap range\ xn \neq \{\}$
using $xn(5)$ **that** **by**($auto\ simp: md.mdense-def$)
have $infin: infinite (md.mball\ x\ e)$
using $md.perfect-set-mball-infinite[OF\ perfect]$ **that** **by** $simp$
then obtain y **where** $y: y \in md.mball\ x\ e\ y \notin range\ xn$
using $h-fin$ **by**($metis\ inf.absorb-iff2\ inf-commute\ subsetI$)
define e' **where** $e' = Min\ \{d\ y\ xk\ |xk. xk \in md.mball\ x\ e \cap range\ xn\}$
have $fin: finite\ \{d\ y\ xk\ |xk. xk \in md.mball\ x\ e \cap range\ xn\}$
using $finite-imageI[OF\ h-fin,of\ d\ y]$ **by** ($metis\ Setcompr-eq-image$)
have $nen: \{d\ y\ xk\ |xk. xk \in md.mball\ x\ e \cap range\ xn\} \neq \{\}$
using $h-nen$ **by** $auto$
have $e' > 0$
unfolding $e'-def\ Min-gr-iff[OF\ fin\ nen]$
proof safe
fix l
assume $xn\ l \in md.mball\ x\ e$
with y
show $0 < d\ y\ (xn\ l)$
by $auto$
qed
obtain e'' **where** $e'': e'' > 0\ md.mball\ y\ e'' \subseteq md.mball\ x\ e\ y \in md.mball\ y$
 e'' **by** ($meson\ md.centre-in-mball-iff\ md.in-mball\ md.openin-mball\ md.openin-mtopology$
 $y(1)$)
define ε **where** $\varepsilon \equiv min\ e'\ e''$
have $\varepsilon > 0$
using $e''(1)\ \langle e' > 0 \rangle$ **by**($simp\ add: \varepsilon-def$)
then obtain m **where** $m: d\ y\ (xn\ m) < \varepsilon$
using $md.mdense-def2[of\ range\ xn]\ xn(5)\ y(1)$ **by** $fastforce$
consider $xn\ m \in md.mball\ x\ e\ | xn\ m \in P - md.mball\ x\ e$
using $xn(4)$ **by** $auto$
then show $False$
proof cases
case 1
then have $e' \leq d\ y\ (xn\ m)$
using $Min-le-iff[OF\ fin\ nen]$ **by**($auto\ simp: e'-def$)
thus $?thesis$
using m **by**($simp\ add: \varepsilon-def$)
next
case 2
then have $xn\ m \notin md.mball\ y\ e''$
using $e''(2)$ **by** $auto$
hence $e'' \leq d\ y\ (xn\ m)$

```

    using y e'' xn by auto
  thus ?thesis
    using m by(simp add: ε-def)
qed
qed
have hinfin:infinite (md.mball x e ∩ (xn ' {l<..})) if x ∈ P e > 0 for x e l
proof
  assume finite (md.mball x e ∩ xn ' {l<..})
  moreover have finite (md.mball x e ∩ xn ' {..l}) by simp
  moreover have (md.mball x e ∩ (range xn)) = (md.mball x e ∩ xn ' {l<..})
  ∪ (md.mball x e ∩ xn ' {..l})
    by fastforce
  ultimately have finite (md.mball x e ∩ (range xn))
    by auto
  with hinfin'[OF that] show False ..
qed
have infinite (md.mball (xn n) ((1/2) ^ Suc n'))
  using md.perfect-set-mball-infinite[OF perfect] xn(4)[of n] by simp
then obtain x where x: x ∈ md.mball (xn n) ((1/2) ^ Suc n') x ≠ xn n
  by (metis finite-insert finite-subset infinite-imp-nonempty singletonI subsetI)
then obtain e where e: e > 0 md.mball x e ⊆ md.mball (xn n) ((1/2) ^ Suc
n') x ∈ md.mball x e
  by (meson md.centre-in-mball-iff md.in-mball md.openin-mball md.openin-mtopology)
have d (xn n) x > 0
  using xn x by simp
then obtain m' where m': m' - 1 > 0 (1/2) ^ (m' - 1) < d (xn n) x
  by (metis One-nat-def diff-Suc-Suc diff-zero one-less-numeral-iff reals-power-lt-ex
semiring-norm(76))
define m where m ≡ max m' (max n' (Suc n))
then have m ≥ m' m ≥ n' m ≥ Suc n by simp-all
hence m: m - 1 > 0 (1/2) ^ (m - 1) < d (xn n) x m > n
  using m' less-trans[OF - m'(2),of (1 / 2) ^ (m - 1)]
  by auto (metis diff-less-mono le-eq-less-or-eq)
define ε where ε ≡ min e (d (xn n) x - (1/2) ^ (m - 1))
have ε > 0
  using e m by(simp add: ε-def)
have ball-le:md.mball x ε ⊆ md.mball (xn n) ((1 / 2) ^ Suc n')
  using e(2) by(auto simp add: ε-def)
obtain k where k: xn k ∈ md.mball x ε k > m
  using ⟨ε > 0⟩ infinite-imp-nonempty[OF hinfin,of - ε] x(1) by fastforce
show ∃ m > n. n' < m ∧ (1 / 2) ^ (m - 1) < d (xn n) (xn m) ∧ d (xn n) (xn
m) < (1 / 2) ^ Suc n'
proof(intro exI[where x=k] conjI)
  have (1 / 2) ^ (k - 1) < (1 / (2 :: real)) ^ (m - 1)
    using k(2) m(3) by simp
  also have ... = d (xn n) x + ((1/2) ^ (m - 1) - d (xn n) x) by simp
  also have ... < d (xn n) x - d (xn k) x
    using k by(auto simp: ε-def md.commute)
  also have ... ≤ d (xn n) (xn k)

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      using xn x md.md-dist-reverse-triangle[of xn n x xn k] by(auto simp:
md commute)
      finally show (1 / 2) ^ (k - 1) < d (xn n) (xn k) .
      qed(use <m ≥ n'> k ball-le m(3) in auto)
    qed
    have jn n > n ∧ md.mcball (xn (jn n)) ((1/2)^(jn n)) ⊆ md.mball (xn n)
((1/2)^(jn n)) - md.mball (xn n) ((1/2)^(jn n)) for n
      unfolding jn-def
    proof(rule LeastI-ex)
      obtain m where m:m > n (1 / 2) ^ (m - 1) < d (xn n) (xn m) d (xn n)
(xn m) < (1 / 2) ^ Suc n
      using dxm:xn by auto
      show ∃x>n. md.mcball (xn x) ((1 / 2) ^ x) ⊆ md.mball (xn n) ((1 / 2) ^ n)
- md.mball (xn n) ((1 / 2) ^ x)
      proof(safe intro!: exI[where x=m] m(1))
        fix x
        assume h:x ∈ md.mcball (xn m) ((1 / 2) ^ m)
        have 1:d (xn n) x < (1 / 2) ^ n
        proof -
          have d (xn n) x < (1 / 2) ^ Suc n + (1 / 2) ^ m
            using m(3) md.triangle[OF xn(4)[of n] xn(4)[of m],of x] h by auto
          also have ... ≤ (1 / 2) ^ Suc n + (1 / 2) ^ Suc n
          by (metis Suc-lessI add-mono divide-less-eq-1-pos divide-pos-pos less-eq-real-def
m(1) one-less-numeral-iff power-strict-decreasing-iff semiring-norm(76) zero-less-numeral
zero-less-one)
        finally show ?thesis by simp
      qed
      have 2:(1 / 2) ^ m ≤ d (xn n) x
      proof -
        have (1 / 2) ^ (m - 1) < d (xn n) x + (1 / 2) ^ m
          using order.strict-trans2[OF m(2) md.triangle[OF xn(4)[of n] - xn(4)[of
m]]] h md commute by fastforce
        hence (1 / 2) ^ (m - 1) - (1 / 2) ^ m ≤ d (xn n) x
          by simp
        thus ?thesis
          using not0-implies-Suc[OF gr-implies-not0[OF m(1)]] by auto
      qed
      show x ∈ md.mball (xn n) ((1 / 2) ^ n)
        x ∈ md.mball (xn n) ((1 / 2) ^ m) ⇒ False
      using xn h 1 2 by auto
    qed
  qed
  hence jn: ∧n. jn n > n ∧n. md.mcball (xn (jn n)) ((1/2)^(jn n)) ⊆ md.mball
(xn n) ((1/2)^(jn n)) - md.mball (xn n) ((1/2)^(jn n))
    by simp-all
  have kn n > jn n ∧ md.mcball (xn (kn n)) ((1/2)^(kn n)) ⊆ md.mball (xn n)
((1/2)^(jn n)) for n
    unfolding kn-def
  proof(rule LeastI-ex)

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obtain  $m$  where  $m:m > jn\ n\ d\ (xn\ n)\ (xn\ m) < (1 / 2) \wedge Suc\ (jn\ n)$ 
using  $d\ x\ m\ x\ n$  by  $blast$ 
show  $\exists x > jn\ n.\ md.mcball\ (xn\ x)\ ((1 / 2) \wedge x) \subseteq md.mball\ (xn\ n)\ ((1 / 2) \wedge$ 
 $jn\ n)$ 
proof ( $intro\ exI$  [where  $x=m$ ]  $conjI$ )
show  $md.mcball\ (xn\ m)\ ((1 / 2) \wedge m) \subseteq md.mball\ (xn\ n)\ ((1 / 2) \wedge jn\ n)$ 
proof
fix  $x$ 
assume  $h: x \in md.mcball\ (xn\ m)\ ((1 / 2) \wedge m)$ 
have  $d\ (xn\ n)\ x < (1 / 2) \wedge Suc\ (jn\ n) + (1 / 2) \wedge m$ 
using  $md.triangle$  [ $OF\ xn(4)$ ] [ $of\ n$ ] [ $xn(4)$ ] [ $of\ m$ ]  $h\ m(2)$  by  $fastforce$ 
also have  $\dots \leq (1 / 2) \wedge Suc\ (jn\ n) + (1 / 2) \wedge Suc\ (jn\ n)$ 
by ( $metis\ Suc-le-eq\ add-mono\ dual-order.refl\ less-divide-eq-1-pos\ linorder-not-less$ 
 $m(1)\ not-numeral-less-one\ power-decreasing\ zero-le-divide-1-iff\ zero-le-numeral\ zero-less-numeral$ )
finally show  $x \in md.mball\ (xn\ n)\ ((1 / 2) \wedge jn\ n)$ 
using  $xn(4)\ h$  by  $auto$ 
qed
qed ( $use\ m(1)$  in  $auto$ )
qed
hence  $kn: \bigwedge n.\ kn\ n > jn\ n \bigwedge n.\ md.mcball\ (xn\ (kn\ n))\ ((1/2) \wedge (kn\ n)) \subseteq$ 
 $md.mball\ (xn\ n)\ ((1/2) \wedge (jn\ n))$ 
by  $simp-all$ 
have  $jnkn-pos: jn\ n > 0\ kn\ n > 0$  for  $n$ 
using  $not0-implies-Suc$  [ $OF\ gr-implies-not0$ ] [ $OF\ jn(1)$ ] [ $of\ n$ ]]  $kn(1)$  [ $of\ n$ ] by  $auto$ 

define  $bn :: real\ list \Rightarrow nat$ 
where  $bn \equiv rec-list\ 1\ (\lambda a\ l\ t.\ if\ a = 0\ then\ jn\ t\ else\ kn\ t)$ 
have  $bn-simp: bn\ [] = 1\ bn\ (a\ \#\ l) = (if\ a = 0\ then\ jn\ (bn\ l)\ else\ kn\ (bn\ l))$  for
 $a\ l$ 
by ( $simp-all\ add: bn-def$ )
define  $to-listn :: (nat \Rightarrow real) \Rightarrow nat \Rightarrow real\ list$ 
where  $to-listn \equiv (\lambda x.\ rec-nat\ []\ (\lambda n\ t.\ x\ n\ \# t))$ 
have  $to-listn-simp: to-listn\ x\ 0 = []\ to-listn\ x\ (Suc\ n) = x\ n\ \# to-listn\ x\ n$  for  $x$ 
 $n$ 
by ( $simp-all\ add: to-listn-def$ )
have  $to-listn-eq: (\bigwedge m.\ m < n \Longrightarrow x\ m = y\ m) \Longrightarrow to-listn\ x\ n = to-listn\ y\ n$ 
for  $x\ y\ n$ 
by ( $induction\ n$ ) ( $auto\ simp: to-listn-simp$ )
have  $bn-gtn: bn\ (to-listn\ x\ n) > n$  for  $x\ n$ 
apply ( $induction\ n\ arbitrary: x$ )
using  $jn(1)\ kn(1)$  by ( $auto\ simp: bn-simp\ to-listn-simp$ ) ( $meson\ Suc-le-eq\ le-less$ 
 $less-trans-Suc$ ) +
define  $rn$  where  $rn \equiv (\lambda n.\ Min\ (range\ (\lambda x.\ (1 / 2 :: real) \wedge bn\ (to-listn\ x\ n))))$ 
have  $rn-fin: finite\ (range\ (\lambda x.\ (1 / 2 :: real) \wedge bn\ (to-listn\ x\ n)))$  for  $n$ 
proof -
have  $finite\ (range\ (\lambda x.\ bn\ (to-listn\ x\ n)))$ 
proof ( $induction\ n$ )
case  $ih: (Suc\ n)$ 
have  $(range\ (\lambda x.\ bn\ (to-listn\ x\ (Suc\ n)))) \subseteq (range\ (\lambda x.\ jn\ (bn\ (to-listn\ x$ 

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n)))) ∪ (range (λx. kn (bn (to-listn x n))))
  by(auto simp: to-listn-simp bn-simp)
  moreover have finite ...
  using ih finite-range-imageI by auto
  ultimately show ?case by(rule finite-subset)
qed(simp add: to-listn-simp)
thus ?thesis
  using finite-range-imageI by blast
qed
have rn-nen: (range (λx. (1 / 2 :: real) ^ bn (to-listn x n))) ≠ {} for n
  by simp
have rn-pos: 0 < rn n for n
  by(simp add: Min-gr-iff[OF rn-fin rn-nen] rn-def)
have rn-less: rn n < (1/2) ^ n for n
  using bn-gtn[of n] by(auto simp: rn-def Min-less-iff[OF rn-fin rn-nen])
have cball-le-ball: md.mcball (xn (bn (a#l))) ((1/2) ^ (bn (a#l))) ⊆ md.mball (xn
(bn l)) ((1/2) ^ (bn l)) for a l
  using kn(2)[of bn l] less-imp-le[OF jn(1)] jn(2) md.mball-subset-concentric[of
(1 / 2) ^ jn (bn l) (1 / 2) ^ bn l xn (bn l)]
  by(auto simp: bn-simp)
hence cball-le: md.mcball (xn (bn (a#l))) ((1/2) ^ (bn (a#l))) ⊆ md.mcball (xn
(bn l)) ((1/2) ^ (bn l)) for a l
  using md.mball-subset-mcball by blast
have cball-disj: md.mcball (xn (bn (0#l))) ((1/2) ^ (bn (0#l))) ∩ md.mcball (xn
(bn (1#l))) ((1/2) ^ (bn (1#l))) = {} for l
  using jn(2) kn(2) by(auto simp: bn-simp)
have ∀x. ∃l. l ∈ P ∧ (∩ n. md.mcball (xn (bn (to-listn x n))) ((1 / 2) ^ bn
(to-listn x n))) = {l}
  proof
  fix x
  show ∃l. l ∈ P ∧ (∩ n. md.mcball (xn (bn (to-listn x n))) ((1 / 2) ^ bn
(to-listn x n))) = {l}
  proof(safe intro!: md.mcomplete-nest-sing[THEN iffD1, OF mdc, rule-format])
  show md.mcball (xn (bn (to-listn x n))) ((1 / 2) ^ bn (to-listn x n)) = {}
⇒ False for n
  using md.mcball-eq-empty xn(4) by auto
  next
  show decseq (λn. md.mcball (xn (bn (to-listn x n))) ((1 / 2) ^ bn (to-listn x
n)))
  by(intro decseq-SucI, simp add: to-listn-simp cball-le)
  next
  fix e :: real
  assume 0 < e
  then obtain N where N: (1 / 2) ^ N < e
  by (meson reals-power-lt-ex rel-simps(49) rel-simps(9))
  show ∃ n a. md.mcball (xn (bn (to-listn x n))) ((1 / 2) ^ bn (to-listn x n))
⊆ md.mcball a e
  proof(safe intro!: exI[where x=N] exI[where x=xn (bn (to-listn x N))])
  fix y

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    assume  $y \in \text{md.mcball} (xn (bn (to-listn x N))) ((1 / 2) \wedge bn (to-listn x N))$ 
  then have  $y \in \text{md.mcball} (xn (bn (to-listn x N))) ((1 / 2) \wedge N)$ 
    using  $\text{md.mcball-subset-concentric}[OF \text{power-decreasing}[OF \text{less-imp-le}[OF \text{bn-gtn}[of N x]],of 1/2]]$ 
    by fastforce
  thus  $y \in \text{md.mcball} (xn (bn (to-listn x N))) e$ 
    using  $N \langle 0 < e \rangle$  by auto
  qed
  qed
  then obtain  $f$  where  $f: \bigwedge x. f x \in P \wedge x. (\bigcap n. \text{md.mcball} (xn (bn (to-listn x n)))) ((1 / 2) \wedge bn (to-listn x n)) = \{f x\}$ 
    by metis
  hence  $f': \bigwedge n x. f x \in \text{md.mcball} (xn (bn (to-listn x n))) ((1 / 2) \wedge bn (to-listn x n))$ 
    by blast
  have  $f'': f x \in \text{md.mball} (xn (bn (to-listn x n))) ((1 / 2) \wedge bn (to-listn x n))$  for  $n x$ 
    using  $f'[of x \text{Suc } n] \text{cball-le-ball}[of - to-listn x n]$  by (fastforce simp: to-listn-simp)
  interpret  $\text{bdmd}: \text{Submetric } P d f' (\prod_E i \in \text{UNIV}. \{0,1\})$ 
    by standard (use f in auto)
  have  $\text{bdmd-sub}: \text{bdmd.sub.mtopology} = \text{subtopology } X (f' (\prod_E i \in \text{UNIV}. \{0,1\}))$ 
    using  $f(1) \text{Int-absorb1}[of f' (UNIV \rightarrow_E \{0,1\}) P]$  by (fastforce simp: bdmd.mtopology-submetric d subtopology-subtopology)
  let  $?d = \lambda x y. \text{if } (x = 0 \vee x = 1) \wedge (y = 0 \vee y = 1) \text{ then } \text{dist } x y \text{ else } 0$ 
  interpret  $d01: \text{Metric-space } \{0,1::\text{real}\} ?d$ 
    by (auto simp: Metric-space-def)
  have  $d01: d01.mtopology = \text{top-of-set } \{0,1\}$ 
  proof -
    have  $d01.mtopology = \text{Metric-space.mtopology } \{0,1\} \text{ dist}$ 
    by (auto intro!: Metric-space-eq-mtopology simp: Metric-space-def metric-space-class.dist-commute)
    also have  $\dots = \text{top-of-set } \{0,1\}$ 
    by (auto intro!: Submetric.mtopology-submetric[of UNIV dist \{0,1::real\},simplified])
  simp: Submetric-def Metric-space-def Submetric-axioms-def dist-real-def
  finally show  $?thesis$  .
  qed
  interpret  $\text{pd}: \text{Product-metric } 1/2 \text{ UNIV id id } \lambda-. \{0,1::\text{real}\} \lambda-. ?d 1$ 
    by (auto intro!: product-metric-natI d01.Metric-space-axioms)
  have  $\text{mpd-top}: \text{pd.Product-metric.mtopology} = \text{Cantor-space-topology}$ 
    by (auto simp: pd.Product-metric-mtopology-eq[symmetric] d01 intro!: product-topology-cong)
  define  $\text{def-at}$  where  $\text{def-at } x y \equiv \text{LEAST } n. x n \neq y n$  for  $x y :: \text{nat} \Rightarrow \text{real}$ 
  have  $\text{def-atxy}: \bigwedge n. n < \text{def-at } x y \implies x n = y n \wedge x (\text{def-at } x y) \neq y (\text{def-at } x y)$ 
  if  $x \neq y$  for  $x y$ 
  proof -
    have  $\exists n. x n \neq y n$ 
    using that by auto
    from LeastI-ex[OF this]

```


show $\bigwedge n. n < \text{def-at } x \ y \implies x \ n = y \ n \ x \ (\text{def-at } x \ y) \neq y \ (\text{def-at } x \ y)$
using *not-less-Least* **by**(*auto simp: def-at-def*)
qed
have *def-at-le-if*: $\text{pd.product-dist } x \ y \leq (1/2)^{\wedge n} \implies n \leq \text{def-at } x \ y$ **if** *assm*: $x \neq y$
 $x \in (\prod_E i \in \text{UNIV}. \{0,1\}) \ y \in (\prod_E i \in \text{UNIV}. \{0,1\})$ **for** $x \ y \ n$
proof –
assume *h*: $\text{pd.product-dist } x \ y \leq (1 / 2)^{\wedge n}$
have $x \ m = y \ m$ **if** *m-less-n*: $m < n$ **for** m
proof(*rule ccontr*)
assume *nen*: $x \ m \neq y \ m$
then have $?d \ (x \ m) \ (y \ m) = 1$
using *assm(2,3)* **by**(*auto simp: submetric-def*)
hence $1 \leq 2^{\wedge m} * \text{pd.product-dist } x \ y$
using *pd.product-dist-geq[of m m,simplified,OF assm(2,3)]* **by** *simp*
hence $(1/2)^{\wedge m} \leq 2^{\wedge m} * (1/2)^{\wedge m} * \text{pd.product-dist } x \ y$ **by** *simp*
hence $(1/2)^{\wedge m} \leq \text{pd.product-dist } x \ y$ **by** (*simp add: power-one-over*)
also have $\dots \leq (1 / 2)^{\wedge n}$
by(*simp add: h*)
finally show *False*
using *that* **by** *auto*
qed
thus $n \leq \text{def-at } x \ y$
by (*meson def-atxy(2) linorder-not-le that(1)*)
qed
have *def-at-le-then*: $\text{pd.product-dist } x \ y \leq 2 * (1/2)^{\wedge n}$ **if** *assm*: $x \neq y$ $x \in (\prod_E$
 $i \in \text{UNIV}. \{0,1\}) \ y \in (\prod_E i \in \text{UNIV}. \{0,1\}) \ n \leq \text{def-at } x \ y$ **for** $x \ y \ n$
proof –
have $\bigwedge m. m < n \implies x \ m = y \ m$
by (*metis def-atxy(1) order-less-le-trans that(4)*)
hence $1: \bigwedge m. m < n \implies ?d \ (x \ m) \ (y \ m) = 0$
by (*simp add: submetric-def*)
have $\text{pd.product-dist } x \ y = (\sum i. (1/2)^{\wedge(i+n)} * (?d \ (x \ (i+n)) \ (y \ (i+n))))$
 $+ (\sum i < n. (1/2)^{\wedge i} * (?d \ (x \ i) \ (y \ i)))$
using *assm* *pd.product-dist-summable'[simplified]* **unfolding** *product-dist-def*
id-apply **by**(*auto intro!: suminf-split-initial-segment simp: product-dist-def*)
also have $\dots = (\sum i. (1/2)^{\wedge(i+n)} * (?d \ (x \ (i+n)) \ (y \ (i+n))))$
by(*simp add: 1*)
also have $\dots \leq (\sum i. (1/2)^{\wedge(i+n)})$
using *pd.product-dist-summable'* **unfolding** *id-apply* **by**(*auto intro!: suminf-le*
summable-ignore-initial-segment)
finally show *?thesis*
using *pd.nsum-of-rK[of n]* **by** *simp*
qed
have *d-le-def*: $d \ (f \ x) \ (f \ y) \leq (1/2)^{\wedge(\text{def-at } x \ y)}$ **if** *assm*: $x \neq y$ $x \in (\prod_E i \in \text{UNIV}.$
 $\{0,1\}) \ y \in (\prod_E i \in \text{UNIV}. \{0,1\})$ **for** $x \ y$
proof –
have $1: \text{to-listn } x \ n = \text{to-listn } y \ n$ **if** $n \leq \text{def-at } x \ y$ **for** n
proof –
have $\bigwedge m. m < n \implies x \ m = y \ m$

```

    by (metis def-atxy(1) order-less-le-trans that)
  then show ?thesis
    by(auto intro!: to-listn-eq)
  qed
  have f x ∈ md.mcball (xn (bn (to-listn x (def-at x y)))) ((1 / 2) ^ bn (to-listn
x (def-at x y)))
    f y ∈ md.mcball (xn (bn (to-listn x (def-at x y)))) ((1 / 2) ^ bn (to-listn x
(def-at x y)))
    using f'[of x def-at x y] f'[of y def-at x y] by(auto simp: 1[OF order-refl])
  hence d (f x) (f y) ≤ 2 * (1 / 2) ^ bn (to-listn x (def-at x y))
  using f(1) by(auto intro!: md.mdiameter-is-sup'[OF - - md.mdiameter-cball-leq])
  also have ... ≤ (1/2) ^ (def-at x y)
  proof -
    have Suc (def-at x y) ≤ bn (to-listn x (def-at x y))
      using bn-gtn[of def-at x y x] by simp
    hence (1 / 2) ^ bn (to-listn x (def-at x y)) ≤ (1 / 2 :: real) ^ Suc (def-at x
y)
      using power-decreasing-iff[OF pd.r] by blast
    thus ?thesis
      by simp
  qed
  finally show d (f x) (f y) ≤ (1/2) ^ (def-at x y) .
  qed
  have fy-in:f y ∈ md.mcball (xn (bn (to-listn x m))) ((1/2) ^ bn (to-listn x m))
⇒ ∀ l < m. x l = y l if assm:x ∈ (ΠE i ∈ UNIV. {0,1}) y ∈ (ΠE i ∈ UNIV. {0,1})
for x y m
  proof(induction m)
    case ih:(Suc m)
    have f y ∈ md.mcball (xn (bn (to-listn x m))) ((1 / 2) ^ bn (to-listn x m))
      using ih(2) cball-le by(fastforce simp: to-listn-simp)
    with ih(1) have k:k < m ⇒ x k = y k for k by simp
    show ?case
  proof safe
    fix l
    assume l < Suc m
    then consider l < m | l = m
      using ‹l < Suc m› by fastforce
    thus x l = y l
  proof cases
    case 2
    have 3:f y ∈ md.mcball (xn (bn (y l # to-listn y l))) ((1 / 2) ^ bn (y l #
to-listn y l))
      using f'[of y Suc l] by(simp add: to-listn-simp)
    have 4:f y ∈ md.mcball (xn (bn (x l # to-listn y l))) ((1 / 2) ^ bn (x l #
to-listn y l))
      using ih(2) to-listn-eq[of m x y,OF k] by(simp add: to-listn-simp 2)
    show ?thesis
  proof(rule ccontr)
    assume x l ≠ y l

```

```

then consider  $x\ l = 0\ y\ l = 1 \mid x\ l = 1\ y\ l = 0$ 
  using  $asm(1,2)$  by( $auto\ simp: PiE-def\ Pi-def$ )  $metis$ 
  thus  $False$ 
  by cases ( $use\ cball-disj[of\ to-listn\ y\ l]\ 3\ 4$  in  $auto$ )
qed
qed( $simp\ add: k$ )
qed
qed  $simp$ 
have  $d-le-rn-then: \exists e>0. \forall y \in (\prod_E i \in UNIV. \{0,1\}). x \neq y \longrightarrow d(f\ x)\ (f\ y) < e \longrightarrow n \leq def-at\ x\ y$  if  $asm: x \in (\prod_E i \in UNIV. \{0,1\})$  for  $x\ n$ 
proof( $safe\ intro!: exI[where\ x=(1/2)^{\wedge}bn\ (to-listn\ x\ n) - d(xn\ (bn\ (to-listn\ x\ n)))\ (f\ x)]$ )
  show  $0 < (1 / 2)^{\wedge}bn\ (to-listn\ x\ n) - d(xn\ (bn\ (to-listn\ x\ n)))\ (f\ x)$ 
  using  $f''$  by  $auto$ 
next
fix  $y$ 
assume  $h: y \in (\prod_E i \in UNIV. \{0,1\})\ d(f\ x)\ (f\ y) < (1 / 2)^{\wedge}bn\ (to-listn\ x\ n) - d(xn\ (bn\ (to-listn\ x\ n)))\ (f\ x)\ x \neq y$ 
then have  $f\ y \in md.mcball(xn\ (bn\ (to-listn\ x\ n)))\ ((1/2)^{\wedge}bn\ (to-listn\ x\ n))$ 
  using  $md.triangle[OF\ xn(4)[of\ bn\ (to-listn\ x\ n)]\ f(1)[of\ x]\ f(1)[of\ y]$ 
  by( $simp\ add: xn(4)[of\ bn\ (to-listn\ x\ n)]\ f(1)[of\ y]\ md.mcball-def$ )
with  $fy-in[OF\ asm\ h(1)]$  have  $\forall m < n. x\ m = y\ m$ 
  by  $auto$ 
thus  $n \leq def-at\ x\ y$ 
  by ( $meson\ def-atxy(2)\ linorder-not-le\ h(3)$ )
qed
have  $0: f'(\prod_E i \in UNIV. \{0,1\}) \subseteq topspace\ X$ 
  using  $f(1)\ P(2)$  by  $auto$ 
have  $1: continuous-map\ pd.Product-metric.mtopology\ bmd.sub.mtopology\ f$ 
unfolding  $pd.Product-metric.metric-continuous-map[OF\ bmd.sub.Metric-space-axioms]$ 
proof  $safe$ 
  fix  $x :: nat \Rightarrow real$  and  $\varepsilon :: real$ 
  assume  $h: x \in (\prod_E i \in UNIV. \{0,1\})\ 0 < \varepsilon$ 
  then obtain  $n$  where  $n: (1/2)^{\wedge}n < \varepsilon$ 
  using  $real-arch-pow-inv[OF\ -\ pd.r(2)]$  by  $auto$ 
  show  $\exists \delta > 0. \forall y. y \in UNIV \rightarrow_E \{0, 1\} \wedge pd.product-dist\ x\ y < \delta \longrightarrow d(f\ x)\ (f\ y) < \varepsilon$ 
proof( $safe\ intro!: exI[where\ x=(1/2)^{\wedge}n]$ )
  fix  $y$ 
  assume  $y: y \in (\prod_E i \in UNIV. \{0,1\})\ pd.product-dist\ x\ y < (1 / 2)^{\wedge}n$ 
  consider  $x = y \mid x \neq y$  by  $auto$ 
  thus  $d(f\ x)\ (f\ y) < \varepsilon$ 
proof  $cases$ 
  case  $1$ 
  with  $y(1)\ h\ md.zero[OF\ f(1)[of\ y]\ f(1)[of\ y]$ 
  show  $?thesis$  by  $simp$ 
next
  case  $2$ 
  then have  $n \leq def-at\ x\ y$ 

```

```

    using h(1) y by(auto intro!: def-at-le-if)
  have d (f x) (f y) ≤ (1/2)^(def-at x y)
    using h(1) y(1) by(auto simp: d-le-def[OF 2 h(1) y(1)])
  also have ... ≤ (1/2)^(n)
    using ⟨n ≤ def-at x y⟩ by simp
  finally show ?thesis
    using n by simp
qed
qed simp
qed
have 2: open-map pd.Product-metric.mtopology bmd.sub.mtopology f
proof -
  have open-map (mtopology-of pd.Product-metric.Self) (subtopology (mtopology-of
md.Self) (f ' mspace pd.Product-metric.Self)) f
  proof(safe intro!: Metric-space-open-map-from-dist)
    fix x :: nat ⇒ real and ε :: real
    assume h:x ∈ mspace pd.Product-metric.Self 0 < ε
    then have x:x ∈ (ΠE i ∈ UNIV. {0,1}) by simp
    from h obtain n where n: (1/2)^(n) < ε
      using real-arch-pow-inv[OF - pd.r(2)] by auto
    obtain e where e: e > 0 ∧ y. y ∈ (ΠE i ∈ UNIV. {0,1}) ⇒ x ≠ y ⇒ d (f
x) (f y) < e ⇒ Suc n ≤ def-at x y
      using d-le-rn-then[OF x,of Suc n] by auto
    show ∃ δ > 0. ∀ y ∈ mspace pd.Product-metric.Self. mdist md.Self (f x) (f y) <
δ → mdist pd.Product-metric.Self x y < ε
      unfolding md.mdist-Self pd.Product-metric.mspace-Self pd.Product-metric.mdist-Self
      proof(safe intro!: exI[where x=e])
        fix y
        assume y:y ∈ (ΠE i ∈ UNIV. {0,1}) and d (f x) (f y) < e
        then have d':d (f x) (f y) < e
          using h(1) by simp
        consider x = y | x ≠ y by auto
        thus pd.product-dist x y < ε
          by cases (use pd.Product-metric.zero[OF y y] h(2) def-at-le-then[OF - x y
e(2)[OF y - d']] n in auto)
        qed(use e(1) in auto)
      qed(use f in auto)
    thus ?thesis
      by (simp add: bmd.mtopology-submetric mtopology-of-def)
  qed
have 3: f '(topspace pd.Product-metric.mtopology) = topspace bmd.sub.mtopology
  by simp
have 4: inj-on f (topspace pd.Product-metric.mtopology)
  unfolding pd.Product-metric.topspace-mtopology
proof
  fix x y
  assume h:x ∈ (ΠE i ∈ UNIV. {0,1}) y ∈ (ΠE i ∈ UNIV. {0,1}) f x = f y
  show x = y
  proof

```

```

fix n
  have f y ∈ md.mcball (xn (bn (to-listn x (Suc n)))) ((1/2) ^ bn (to-listn x
(Suc n)))
  using f'[of x Suc n] by(simp add: h)
  thus x n = y n
  using fy-in[OF h(1,2),of Suc n] by simp
qed
qed
show ?thesis
  using homeomorphic-map-imp-homeomorphic-space[OF bijective-open-imp-homeomorphic-map[OF
1 2 3 4]] 0
  by (metis (no-types, lifting) assms(1) bdmd-sub-completely-metrizable-space-homeo-image-gdelta-in
mpd-top Polish-space-Cantor-space Polish-space-def)
qed

```

2.4 Borel Spaces generated from Polish Spaces

lemma *closedin-clopen-topology*:

```

assumes Polish-space X closedin X a
shows ∃ X'. Polish-space X' ∧ (∀ u. openin X u → openin X' u) ∧ topspace X
= topspace X' ∧ sets (borel-of X) = sets (borel-of X') ∧ openin X' a ∧ closedin
X' a
proof –
  have p1:Polish-space (subtopology X a)
    by (simp add: assms Polish-space-closedin)
  then obtain da' where da': Metric-space a da' subtopology X a = Metric-space.mtopology
a da' Metric-space.mcomplete a da'
    by (metis Metric-space.topspace-mtopology assms(2) closedin-subset-completely-metrizable-space-def
Polish-space-imp-completely-metrizable-space topspace-subtopology-subset)
  define da where da = Metric-space.capped-dist da' (1/2)
  have da: subtopology X a = Metric-space.mtopology a da Metric-space.mcomplete
a da
    using da' by(auto simp: da-def Metric-space.mtopology-capped-metric Met-
ric-space.mcomplete-capped-metric)
  interpret pa: Metric-space a da
    using da' by(simp add: Metric-space.capped-dist da-def)
  have da-bounded: ∧x y. da x y < 1
    using da' by(auto simp: da-def Metric-space.capped-dist-def)
  have p2:Polish-space (subtopology X (topspace X – a))
    by (meson assms(1) assms(2) closedin-def Polish-space-openin)
  then obtain db' where db': Metric-space (topspace X – a) db' subtopology X
(topspace X – a) = Metric-space.mtopology (topspace X – a) db' Metric-space.mcomplete
(topspace X – a) db'
    by (metis Diff-subset Metric-space.topspace-mtopology-completely-metrizable-space-def
Polish-space-imp-completely-metrizable-space topspace-subtopology-subset)
  define db where db = Metric-space.capped-dist db' (1/2)
  have db: subtopology X (topspace X – a) = Metric-space.mtopology (topspace X
– a) db Metric-space.mcomplete (topspace X – a) db
    using db' by(auto simp: db-def Metric-space.mtopology-capped-metric Met-

```

```

ric-space.mcomplete-capped-metric)
interpret pb: Metric-space topspace X - a db
  using db' by(simp add: Metric-space.capped-dist db-def)
have db-bounded:  $\bigwedge x y. db\ x\ y < 1$ 
  using db' by(auto simp: db-def Metric-space.capped-dist-def)
interpret p: Sum-metric UNIV  $\lambda b. \text{if } b \text{ then } a \text{ else topspace } X - a$   $\lambda b. \text{if } b \text{ then } da \text{ else } db$ 
  using da db da-bounded db-bounded by(auto intro!: sum-metricI simp: disjoint-family-on-def pa.Metric-space-axioms pb.Metric-space-axioms)
have 0:  $(\bigcup i. \text{if } i \text{ then } a \text{ else topspace } X - a) = \text{topspace } X$ 
  using closedin-subset assms by auto

have 1: sets (borel-of X) = sets (borel-of p.Sum-metric.mtopology)
proof -
  have sigma-sets (topspace X) (Collect (openin X)) = sigma-sets (topspace X)
(Collect (openin p.Sum-metric.mtopology))
  proof(rule sigma-sets-eqI)
    fix a
    assume a  $\in$  Collect (openin X)
    then have openin p.Sum-metric.mtopology a
      by(simp only: p.openin-mtopology-iff) (auto simp: 0 da(1)[symmetric]
db(1)[symmetric] openin-subtopology dest: openin-subset)
    thus a  $\in$  sigma-sets (topspace X) (Collect (openin p.Sum-metric.mtopology))
      by auto
  next
  interpret s: sigma-algebra topspace X sigma-sets (topspace X) (Collect (openin
X))
    by(auto intro!: sigma-algebra-sigma-sets openin-subset)
  fix b
  assume b  $\in$  Collect (openin p.Sum-metric.mtopology)
  then have openin p.Sum-metric.mtopology b by auto
    then have b:  $b \subseteq \text{topspace } X$  openin (subtopology X a) (b  $\cap$  a) openin
(subtopology X (topspace X - a)) (b  $\cap$  (topspace X - a))
    by(simp-all only: p.openin-mtopology-iff, insert 0 da(1) db(1)) (auto simp:
all-bool-eq)
  have [simp]: (b  $\cap$  a)  $\cup$  (b  $\cap$  (topspace X - a)) = b
    using Diff-partition b(1) by blast
  have (b  $\cap$  a)  $\cup$  (b  $\cap$  (topspace X - a))  $\in$  sigma-sets (topspace X) (Collect
(openin X))
  proof(rule sigma-sets-Un)
    have [simp]: a  $\in$  sigma-sets (topspace X) (Collect (openin X))
    proof -
      have topspace X - (topspace X - a)  $\in$  sigma-sets (topspace X) (Collect
(openin X))
      by(rule sigma-sets.Compl) (use assms in auto)
    thus ?thesis
      using double-diff[OF closedin-subset[OF assms(2)]] by simp
  qed
  from b(2,3) obtain T T' where T: openin X T openin X T' and [simp]: b

```

```

 $\cap a = T \cap a \cap b \cap (\text{topspace } X - a) = T' \cap (\text{topspace } X - a)$ 
  by(auto simp: openin-subtopology)
  show  $b \cap a \in \text{sigma-sets } (\text{topspace } X)$  (Collect (openin X))
     $b \cap (\text{topspace } X - a) \in \text{sigma-sets } (\text{topspace } X)$  (Collect (openin X))
  using T assms by auto
qed
thus  $b \in \text{sigma-sets } (\text{topspace } X)$  (Collect (openin X))
  by simp
qed
thus ?thesis
  by(simp only: sets-borel-of p.Sum-metric.topspace-mtopology) (use 0 in auto)
qed
have 2:  $\bigwedge u. \text{openin } X u \implies \text{openin } p.\text{Sum-metric.mtopology } u$ 
  by(simp only: p.openin-mtopology-iff) (auto simp: all-bool-eq da(1)[symmetric]
db(1)[symmetric] openin-subtopology dest: openin-subset)
  have 3:  $\text{openin } p.\text{Sum-metric.mtopology } a$ 
    by(simp only: p.openin-mtopology-iff) (auto simp: all-bool-eq)
  have 4:  $\text{closedin } p.\text{Sum-metric.mtopology } a$ 
    by (metis 0 2 assms(2) closedin-def p.Sum-metric.topspace-mtopology)
  have 5:  $\text{topspace } X = \text{topspace } p.\text{Sum-metric.mtopology}$ 
    by(simp only: p.Sum-metric.topspace-mtopology) (simp only: 0)
  have 6: Polish-space  $p.\text{Sum-metric.mtopology}$ 
    by(rule p.Polish-spaceI,insert da(2) db(2) p1 p2) (auto simp: da(1) db(1)
Polish-space-def)
  show ?thesis
    by(rule exI[where  $x=p.\text{Sum-metric.mtopology}$ ]) (insert 5 2 6, simp only: 1 3
4 ,auto)
qed

```

lemma Polish-space-union-Polish:

```

  fixes  $X :: \text{nat} \Rightarrow 'a \text{ topology}$ 
  assumes  $\bigwedge n. \text{Polish-space } (X n) \wedge n. \text{topspace } (X n) = Xt \wedge x y. x \in Xt \implies y \in Xt \implies x \neq y \implies \exists Ox Oy. (\forall n. \text{openin } (X n) Ox) \wedge (\forall n. \text{openin } (X n) Oy) \wedge x \in Ox \wedge y \in Oy \wedge \text{disjnt } Ox Oy$ 
  defines  $Xun \equiv \text{topology-generated-by } (\bigcup n. \{u. \text{openin } (X n) u\})$ 
  shows Polish-space  $Xun$ 
proof -
  have  $\text{topspace } Xun = Xt$ 
    using assms(2) by(auto simp: Xun-def dest:openin-subset)
  define  $f :: 'a \Rightarrow \text{nat} \Rightarrow 'a$  where  $f \equiv (\lambda x n. x)$ 
  have continuous-map  $Xun$  (product-topology X UNIV)  $f$ 
    by(auto simp: assms(2) topspace Xun f-def continuous-map-componentwise, auto
simp: Xun-def openin-topology-generated-by-iff continuous-map-def assms(2) dest:openin-subset[of
X -,simplified assms(2)] )
    (insert openin-subopen, fastforce intro!: generate-topology-on.Basis)
  hence 1: continuous-map  $Xun$  (subtopology (product-topology X UNIV) ( $f$  ‘
(topspace Xun)))  $f$ 
    by(auto simp: continuous-map-in-subtopology)
  have 2: inj-on  $f$  (topspace Xun)

```

```

    by(auto simp: inj-on-def f-def dest:fun-cong)
    have 3:  $f'(\text{topspace } Xun) = \text{topspace } (\text{subtopology } (\text{product-topology } X \text{ UNIV}) (f'(\text{topspace } Xun)))$ 
    by(auto simp: topsXun assms(2) f-def)
    have 4:  $\text{open-map } Xun (\text{subtopology } (\text{product-topology } X \text{ UNIV}) (f'(\text{topspace } Xun))) f$ 
    proof(safe intro!: open-map-generated-topo[OF - 2[simplified Xun-def],simplified Xun-def[symmetric]])
      fix u n
      assume u:openin (X n) u
      show openin (subtopology (product-topology X UNIV) (f' topspace Xun)) (f' u)
        unfolding openin-subtopology
        proof(safe intro!: exI[where x={ $\lambda i. \text{if } i = n \text{ then } a \text{ else } b \text{ } i \mid a \text{ } b. a \in u \wedge b \in \text{UNIV} \rightarrow Xt$ }])
          show openin (product-topology X UNIV) { $\lambda i. \text{if } i = n \text{ then } a \text{ else } b \text{ } i \mid a \text{ } b. a \in u \wedge b \in \text{UNIV} \rightarrow Xt$ }
          by(auto simp: openin-product-topology-alt u assms(2) openin-topospace[of X -,simplified assms(2)] intro!: exI[where x= $\lambda i. \text{if } i = n \text{ then } u \text{ else } Xt$ ])
            (auto simp: PiE-def Pi-def, metis openin-subset[OF u,simplified assms(2)] in-mono)
          next
          show  $\bigwedge y. y \in u \implies \exists a \text{ } b. f y = (\lambda i. \text{if } i = n \text{ then } a \text{ else } b \text{ } i) \wedge a \in u \wedge b \in \text{UNIV} \rightarrow Xt$ 
          using assms(2) f-def openin-subset u by fastforce
          next
          show  $\bigwedge y. y \in u \implies f y \in f' \text{ topspace } Xun$ 
          using openin-subset[OF u] by(auto simp: assms(2) topsXun)
          next
          show  $\bigwedge x \text{ } a \text{ } b. xa \in \text{topspace } Xun \implies f xa = (\lambda i. \text{if } i = n \text{ then } a \text{ else } b \text{ } i) \implies a \in u \implies b \in \text{UNIV} \rightarrow Xt \implies f xa \in f' u$ 
          using openin-subset[OF u] by(auto simp: topsXun assms(2)) (metis f-def imageI)
        qed
      qed
    have 5:(subtopology (product-topology X UNIV) (f' topspace Xun)) homeomorphic-space Xun
    using homeomorphic-map-imp-homeomorphic-space[OF bijective-open-imp-homeomorphic-map[OF 1 4 3 2]]
    by(simp add: homeomorphic-space-sym[of Xun])
    show ?thesis
    proof(safe intro!: homeomorphic-Polish-space-aux[OF Polish-space-closedin[OF Polish-space-product] 5] assms)
      show closedin (product-topology X UNIV) (f' topspace Xun)
      proof -
        have 1: openin (product-topology X UNIV) ((UNIV  $\rightarrow_E$  Xt) - f' Xt)
        proof(rule openin-subopen[THEN iffD2])
          show  $\forall x \in (\text{UNIV} \rightarrow_E Xt) - f' Xt. \exists T. \text{openin } (\text{product-topology } X \text{ UNIV}) T \wedge x \in T \wedge T \subseteq (\text{UNIV} \rightarrow_E Xt) - f' Xt$ 

```



```

proof safe
  fix  $x$ 
  assume  $x : x \in UNIV \rightarrow_E Xt \ x \notin f \ ' \ Xt$ 
  have  $\exists n. x \ n \neq x \ 0$ 
  proof(rule ccontr)
    assume  $\nexists n. x \ n \neq x \ 0$ 
    then have  $\forall n. x \ n = x \ 0$  by auto
    hence  $x = (\lambda-. x \ 0)$  by auto
    thus False
    using  $x$  by(auto simp: f-def topsXun assms(2))
  qed
  then obtain  $n$  where  $n : n \neq 0 \ x \ n \neq x \ 0$ 
  by metis
  from assms(3)[OF - - this(2)]  $x$ 
  obtain  $On \ O0$  where  $h : \bigwedge n. \text{openin } (X \ n) \ On \ \bigwedge n. \text{openin } (X \ n) \ O0 \ x \ n$ 
 $\in On \ x \ 0 \in O0 \ \text{disjnt } On \ O0$ 
  by fastforce
  have  $\text{openin } (\text{product-topology } X \ UNIV) ((\lambda x. x \ 0) - \ ' \ O0 \cap \ \text{topspace}$ 
(product-topology } X \ UNIV))
  using continuous-map-product-coordinates[of 0 UNIV X] h(2)[of 0] by
blast
  moreover have  $\text{openin } (\text{product-topology } X \ UNIV) ((\lambda x. x \ n) - \ ' \ On \cap$ 
topspace } (product-topology } X \ UNIV))
  using continuous-map-product-coordinates[of n UNIV X] h(1)[of n] by
blast
  ultimately have  $op : \text{openin } (\text{product-topology } X \ UNIV) ((\lambda T. T \ 0)$ 
 $- \ ' \ O0 \cap \ \text{topspace } (\text{product-topology } X \ UNIV) \cap ((\lambda T. T \ n) - \ ' \ On \cap \ \text{topspace}$ 
(product-topology } X \ UNIV)))
  by auto
  have  $xin : x \in (\lambda T. T \ 0) - \ ' \ O0 \cap \ \text{topspace } (\text{product-topology } X \ UNIV) \cap$ 
 $((\lambda T. T \ n) - \ ' \ On \cap \ \text{topspace } (\text{product-topology } X \ UNIV))$ 
  using  $x \ h(3,4)$  by(auto simp: assms(2))
  have  $\text{subset} : (\lambda T. T \ 0) - \ ' \ O0 \cap \ \text{topspace } (\text{product-topology } X \ UNIV) \cap$ 
 $((\lambda T. T \ n) - \ ' \ On \cap \ \text{topspace } (\text{product-topology } X \ UNIV)) \subseteq (UNIV \rightarrow_E Xt) - \ f$ 
 $\ ' \ Xt$ 
  using  $h(5)$  by(auto simp: assms(2) disjnt-def f-def)

  show  $\exists T. \text{openin } (\text{product-topology } X \ UNIV) \ T \ \wedge \ x \in T \ \wedge \ T \subseteq (UNIV$ 
 $\rightarrow_E Xt) - \ f \ ' \ Xt$ 
  by(rule exI[where x=((\lambda x. x \ 0) - \ ' \ O0 \cap \ \text{topspace } (\text{product-topology } X
UNIV)) \cap ((\lambda x. x \ n) - \ ' \ On \cap \ \text{topspace } (\text{product-topology } X \ UNIV))]) (use op xin
subset in auto)
  qed
  qed
  thus ?thesis
  by(auto simp: closedin-def assms(2) topsXun f-def)
  qed
  qed(simp add: f-def)
  qed

```

lemma *sets-clopen-topology*:

assumes *Polish-space* X $a \in \text{sets (borel-of } X)$

shows $\exists X'. \text{Polish-space } X' \wedge (\forall u. \text{openin } X u \longrightarrow \text{openin } X' u) \wedge \text{topspace } X = \text{topspace } X' \wedge \text{sets (borel-of } X) = \text{sets (borel-of } X') \wedge \text{openin } X' a \wedge \text{closedin } X' a$

proof –

have $a \in \text{sigma-sets (topspace } X) \{U. \text{closedin } X U\}$

using *assms* **by**(*simp add: sets-borel-of-closed*)

thus *?thesis*

proof *induction*

case (*Basic* a)

then show *?case*

by(*simp add: assms closedin-clopen-topology*)

next

case *Empty*

with *polish-space-axioms assms* **show** *?case*

by *auto*

next

case (*Compl* a)

then obtain X' **where** $S': \text{Polish-space } X' (\forall u. \text{openin } X u \longrightarrow \text{openin } X' u) \text{topspace } X = \text{topspace } X' \text{sets (borel-of } X) = \text{sets (borel-of } X') \text{openin } X' a \text{closedin } X' a$

by *auto*

from *closedin-clopen-topology[OF S'(1) S'(6)] S'*

show *?case* **by** *auto*

next

case *ih:(Union* a)

then obtain S_i **where** S_i :

$\bigwedge i. \text{Polish-space } (S_i i) \bigwedge u i. \text{openin } X u \implies \text{openin } (S_i i) u \bigwedge i::\text{nat. topspace } X = \text{topspace } (S_i i) \bigwedge i. \text{sets (borel-of } X) = \text{sets (borel-of } (S_i i)) \bigwedge i. \text{openin } (S_i i) (a i) \bigwedge i. \text{closedin } (S_i i) (a i)$

by *metis*

define S_{un} **where** $S_{un} \equiv \text{topology-generated-by } (\bigcup n. \{u. \text{openin } (S_i n) u\})$

have $S_{un}1: \text{Polish-space } S_{un}$

unfolding $S_{un}\text{-def}$

proof(*safe intro!: Polish-space-union-Polish[where Xt=topspace X]*)

fix $x y$

assume $xy: x \in \text{topspace } X y \in \text{topspace } X x \neq y$

then obtain $Ox Oy$ **where** $Oxy: x \in Ox y \in Oy \text{openin } X Ox \text{openin } X Oy$

disjnt $Ox Oy$

using *metrizable-imp-Hausdorff-space[OF Polish-space-imp-metrizable-space[OF assms(1)]]*

by(*simp only: Hausdorff-space-def*) *metis*

show $\exists Ox Oy. (\forall x. \text{openin } (S_i x) Ox) \wedge (\forall x. \text{openin } (S_i x) Oy) \wedge x \in Ox \wedge y \in Oy \wedge \text{disjnt } Ox Oy$

by(*rule exI[where x=Ox],insert Si(2) Oxy, auto intro!: exI[where x=Oy]*)

qed (*use Si in auto*)

have $S_{un}\text{top:topspace } X = \text{topspace } S_{un}$

```

    using Si(3) by(auto simp: Sun-def dest: openin-subset)
  have Sunsets: sets (borel-of X) = sets (borel-of Sun) (is ?lhs = ?rhs)
  proof -
    have ?lhs = sigma-sets (topspace X) (⋃ n. {u. openin (Si n) u})
    proof
      show sets (borel-of X) ⊆ sigma-sets (topspace X) (⋃ n. {u. openin (Si n)
u})
        using Si(2) by(auto simp: sets-borel-of intro!: sigma-sets-mono)
      next
        show sigma-sets (topspace X) (⋃ n. {u. openin (Si n) u}) ⊆ sets (borel-of
X)
          by(simp add: sigma-sets-le-sets-iff[of borel-of X ⋃ n. {u. openin (Si n)
u},simplified space-borel-of]) (use Si(4) sets-borel-of in fastforce)
        qed
      also have ... = ?rhs
        using borel-of-second-countable'[OF Polish-space-imp-second-countable[OF
Sun1],of ⋃ n. {u. openin (Si n) u}]
        by (simp add: Sun-def Suntop subbase-in-def subset-Pow-Union)
      finally show ?thesis .
    qed
  have Sun-open: ⋀ u i. openin (Si i) u ⇒ openin Sun u
  by(auto simp: Sun-def openin-topology-generated-by-iff intro!: generate-topology-on.Basis)
  have Sun-opena: openin Sun (⋃ i. a i)
  using Sun-open[OF Si(5),simplified Sun-def] by(auto simp: Sun-def openin-topology-generated-by-iff
intro!: generate-topology-on.UN)
  hence closedin Sun (topspace Sun - (⋃ i. a i))
  by auto
  from closedin-clopen-topology[OF Sun1 this]
  show ?case
  using Suntop Sunsets Sun-open[OF Si(2)] Sun-opena
  by (metis closedin-def openin-closedin-eq)
  qed
qed
end

```

3 Standard Borel Spaces

3.1 Standard Borel Spaces

```

theory StandardBorel
  imports Abstract-Metrizable-Topology
begin

  locale standard-borel =
    fixes M :: 'a measure
    assumes Polish-space: ∃ S. Polish-space S ∧ sets M = sets (borel-of S)
begin

```

lemma *singleton-sets*:
assumes $x \in \text{space } M$
shows $\{x\} \in \text{sets } M$
proof –
obtain S **where** $s:\text{Polish-space } S \text{ sets } M = \text{sets } (\text{borel-of } S)$
using *Polish-space* **by** *blast*
have *closedin* $S \{x\}$
using *assms* **by**(*simp add: sets-eq-imp-space-eq*[*OF s(2)*] *closedin-Hausdorff-sing-eq*[*OF metrizable-imp-Hausdorff-space*[*OF Polish-space-imp-metrizable-space*[*OF s(1)*]]] *space-borel-of*)
thus *?thesis*
using *borel-of-closed s* **by** *simp*
qed

corollary *countable-sets*:
assumes $A \subseteq \text{space } M$ *countable* A
shows $A \in \text{sets } M$
using *sets.countable*[*OF singleton-sets assms(2)*] *assms(1)*
by *auto*

lemma *standard-borel-restrict-space*:
assumes $A \in \text{sets } M$
shows *standard-borel* (*restrict-space* $M A$)
proof –
obtain S **where** $s:\text{Polish-space } S \text{ sets } M = \text{sets } (\text{borel-of } S)$
using *Polish-space* **by** *blast*
obtain S' **where** $S':\text{Polish-space } S' \text{ sets } M = \text{sets } (\text{borel-of } S') \text{ openin } S' A$
using *sets-clopen-topology*[*OF s(1),simplified s(2)[symmetric],OF assms*] **by**
auto
show *?thesis*
using *Polish-space-openin*[*OF S'(1,3)*] *S'(2)*
by(*auto simp: standard-borel-def borel-of-subtopology sets-restrict-space intro!*:
exI[**where** $x=\text{subtopology } S' A$])
qed

end

locale *standard-borel-ne = standard-borel +*
assumes *space-ne: space* $M \neq \{\}$
begin

lemma *standard-borel-ne-restrict-space*:
assumes $A \in \text{sets } M$ $A \neq \{\}$
shows *standard-borel-ne* (*restrict-space* $M A$)
using *assms* **by**(*auto simp: standard-borel-ne-def standard-borel-ne-axioms-def standard-borel-restrict-space*)

lemma *standard-borel: standard-borel* M
by(*rule standard-borel-axioms*)

end

lemma *standard-borel-sets*:

assumes *standard-borel M* **and** *sets M = sets N*

shows *standard-borel N*

using *assms* **by**(*simp add: standard-borel-def*)

lemma *standard-borel-ne-sets*:

assumes *standard-borel-ne M* **and** *sets M = sets N*

shows *standard-borel-ne N*

using *assms* **by**(*simp add: standard-borel-def standard-borel-ne-def sets-eq-imp-space-eq*[*OF assms(2)*] *standard-borel-ne-axioms-def*)

lemma *pair-standard-borel*:

assumes *standard-borel M standard-borel N*

shows *standard-borel* ($M \otimes_M N$)

proof –

obtain $S S'$ **where** *hs*:

Polish-space S sets M = sets (borel-of S) Polish-space S' sets N = sets (borel-of S')

using *assms* **by**(*auto simp: standard-borel-def*)

have *sets* ($M \otimes_M N$) = *sets (borel-of (prod-topology S S'))*

unfolding *borel-of-prod*[*OF Polish-space-imp-second-countable*[*OF hs(1)*] *Polish-space-imp-second-countable*[*OF hs(3)*],*symmetric*]

using *sets-pair-measure-cong*[*OF hs(2,4)*] .

thus *?thesis*

unfolding *standard-borel-def* **by**(*auto intro!*: *exI*[**where** $x = \text{prod-topology } S S'$] *simp: Polish-space-prod*[*OF hs(1,3)*])

qed

lemma *pair-standard-borel-ne*:

assumes *standard-borel-ne M standard-borel-ne N*

shows *standard-borel-ne* ($M \otimes_M N$)

using *assms* **by**(*auto simp: pair-standard-borel standard-borel-ne-def standard-borel-ne-axioms-def space-pair-measure*)

lemma *product-standard-borel*:

assumes *countable I*

and $\bigwedge i. i \in I \implies \text{standard-borel } (M i)$

shows *standard-borel* ($\prod_M i \in I. M i$)

proof –

obtain S **where** *hs*:

$\bigwedge i. i \in I \implies \text{Polish-space } (S i) \bigwedge i. i \in I \implies \text{sets } (M i) = \text{sets (borel-of } (S i))$

using *assms(2)* **by**(*auto simp: standard-borel-def*) *metis*

have *sets* ($\prod_M i \in I. M i$) = *sets* ($\prod_M i \in I. \text{borel-of } (S i)$)

using *hs(2)* **by**(*auto intro!*: *sets-PiM-cong*)

also have $\dots = \text{sets (borel-of (product-topology } S I))$

using *assms(1)* *Polish-space-imp-second-countable*[*OF hs(1)*] **by**(*auto intro!*: *sets-PiM-equal-borel-of*)

finally have $1:sets (\prod_{M i \in I. M i}) = sets (borel-of (product-topology S I))$.
show *?thesis*
unfolding *standard-borel-def*
using *assms(1) hs(1)* **by**(*auto intro!*: *exI[where x=product-topology S I] Polish-space-product simp: 1*)
qed

lemma *product-standard-borel-ne*:
assumes *countable I*
and $\bigwedge i. i \in I \implies standard-borel-ne (M i)$
shows *standard-borel-ne* $(\prod_{M i \in I. M i})$
using *assms* **by**(*auto simp: standard-borel-ne-def standard-borel-ne-axioms-def product-standard-borel*)

lemma *closed-set-standard-borel[simp]*:
fixes $U :: 'a :: topological-space set$
assumes *Polish-space (euclidean :: 'a topology) closed U*
shows *standard-borel (restrict-space borel U)*
by(*auto simp: standard-borel-def borel-of-euclidean borel-of-subtopology assms intro!*: *exI[where x=subtopology euclidean U] Polish-space-closedin*)

lemma *closed-set-standard-borel-ne[simp]*:
fixes $U :: 'a :: topological-space set$
assumes *Polish-space (euclidean :: 'a topology) closed U U $\neq \{\}$*
shows *standard-borel-ne (restrict-space borel U)*
using *assms* **by**(*simp add: standard-borel-ne-def standard-borel-ne-axioms-def*)

lemma *open-set-standard-borel[simp]*:
fixes $U :: 'a :: topological-space set$
assumes *Polish-space (euclidean :: 'a topology) open U*
shows *standard-borel (restrict-space borel U)*
by(*auto simp: standard-borel-def borel-of-euclidean borel-of-subtopology assms intro!*: *exI[where x=subtopology euclidean U] Polish-space-openin*)

lemma *open-set-standard-borel-ne[simp]*:
fixes $U :: 'a :: topological-space set$
assumes *Polish-space (euclidean :: 'a topology) open U U $\neq \{\}$*
shows *standard-borel-ne (restrict-space borel U)*
using *assms* **by**(*simp add: standard-borel-ne-def standard-borel-ne-axioms-def*)

lemma *standard-borel-ne-borel[simp]*: *standard-borel-ne (borel :: ('a :: polish-space) measure)*
and *standard-borel-ne-lborel[simp]*: *standard-borel-ne lborel*
unfolding *standard-borel-def standard-borel-ne-def standard-borel-ne-axioms-def*
by(*auto intro!*: *exI[where x=euclidean] simp: borel-of-euclidean*)

lemma *count-space-standard'[simp]*:
assumes *countable I*
shows *standard-borel (count-space I)*

by(rule standard-borel-sets[OF - sets-borel-of-discrete-topology]) (auto simp add: assms Polish-space-discrete-topology standard-borel-def intro!: exI[**where** $x = \text{discrete-topology } I$])

lemma count-space-standard-ne[simp]: standard-borel-ne (count-space (UNIV :: (- :: countable) set))
by (simp add: standard-borel-ne-def standard-borel-ne-axioms-def)

corollary measure-pmf-standard-borel-ne[simp]: standard-borel-ne (measure-pmf (p :: (- :: countable) pmf))
using count-space-standard-ne sets-measure-pmf-count-space standard-borel-ne-sets
by blast

corollary measure-spmf-standard-borel-ne[simp]: standard-borel-ne (measure-spmf (p :: (- :: countable) spmf))
using count-space-standard-ne sets-measure-spmf standard-borel-ne-sets **by** blast

corollary countable-standard-ne[simp]:
standard-borel-ne (borel :: 'a :: {countable,t2-space} measure)
by(simp add: standard-borel-sets[OF - sets-borel-eq-count-space[symmetric]] standard-borel-ne-def standard-borel-ne-axioms-def)

lemma(in standard-borel) countable-discrete-space:
assumes countable (space M)
shows sets M = Pow (space M)
proof safe
fix A
assume $A \subseteq \text{space } M$
with assms **have** countable A
by(simp add: countable-subset)
thus $A \in \text{sets } M$
using $\langle A \subseteq \text{space } M \rangle$ singleton-sets
by(auto intro!: sets.countable[of A])
qed(use sets.sets-into-space in auto)

lemma(in standard-borel) measurable-isomorphic-standard:
assumes M measurable-isomorphic N
shows standard-borel N
proof –
obtain S **where** S:Polish-space S sets M = sets (borel-of S)
using Polish-space **by** auto
from measurable-isomorphic-borels[OF S(2) assms]
obtain S' **where** S': S homeomorphic-space S' \wedge sets N = sets (borel-of S')
by auto
thus ?thesis
by(auto simp: standard-borel-def homeomorphic-Polish-space-aux[OF S(1)] intro!: exI[**where** $x = S'$])
qed

lemma(in *standard-borel-ne*) *measurable-isomorphic-standard-ne*:
assumes *M measurable-isomorphic N*
shows *standard-borel-ne N*
using *measurable-isomorphic-empty2[OF - assms]* **by** (*auto simp: measurable-isomorphic-standard[OF assms]* *standard-borel-ne-def standard-borel-ne-axioms-def space-ne*)

lemma(in *standard-borel*) *standard-borel-embed-measure*:
assumes *inj-on f (space M)*
shows *standard-borel (embed-measure M f)*
using *measurable-embed-measure2'[OF assms]*
by (*auto intro!: measurable-isomorphic-standard exI[where x=f] simp: measurable-isomorphic-def measurable-isomorphic-map-def assms in-sets-embed-measure measurable-def sets.sets-into-space space-embed-measure the-inv-into-into the-inv-into-vimage bij-betw-def*)

corollary(in *standard-borel-ne*) *standard-borel-ne-embed-measure*:
assumes *inj-on f (space M)*
shows *standard-borel-ne (embed-measure M f)*
by (*simp add: assms space-embed-measure space-ne standard-borel-embed-measure standard-borel-ne-axioms-def standard-borel-ne-def*)

lemma
shows *standard-ne-ereal: standard-borel-ne (borel :: ereal measure)*
and *standard-ne-ennreal: standard-borel-ne (borel :: ennreal measure)*
using *Polish-space-ereal Polish-space-ennreal* **by** (*auto simp: standard-borel-ne-def standard-borel-ne-axioms-def standard-borel-def borel-of-euclidean*)

Cantor space \mathcal{C}

definition *Cantor-space* :: (nat \Rightarrow real) measure **where**
Cantor-space \equiv ($\prod_M i \in UNIV. restrict-space borel \{0,1\}$)

lemma *Cantor-space-standard-ne: standard-borel-ne Cantor-space*
by (*auto simp: Cantor-space-def intro!: product-standard-borel-ne*)

lemma *Cantor-space-borel*:
sets (borel-of Cantor-space-topology) = sets Cantor-space
(is ?lhs = -)
proof –
have *?lhs = sets ($\prod_M i \in UNIV. borel-of (top-of-set \{0,1\})$)*
by (*auto intro!: sets-PiM-equal-borel-of[symmetric] second-countable-subtopology*)
thus *?thesis*
by (*simp add: borel-of-subtopology Cantor-space-def borel-of-euclidean*)
qed

Hilbert cube \mathcal{H}

definition *Hilbert-cube* :: (nat \Rightarrow real) measure **where**
Hilbert-cube \equiv ($\prod_M i \in UNIV. restrict-space borel \{0..1\}$)

lemma *Hilbert-cube-standard-ne: standard-borel-ne Hilbert-cube*

by(auto simp: Hilbert-cube-def intro!: product-standard-borel-ne)

lemma Hilbert-cube-borel:

sets (borel-of Hilbert-cube-topology) = sets Hilbert-cube (is ?lhs = -)

proof -

have ?lhs = sets ($\prod_M i \in UNIV. \text{borel-of (top-of-set } \{0..1\})$)

by(auto intro!: sets-PiM-equal-borel-of[symmetric] second-countable-subtopology)

thus ?thesis

by(simp add: borel-of-subtopology Hilbert-cube-def borel-of-euclidean)

qed

3.2 Isomorphism between \mathcal{C} and \mathcal{H}

lemma Cantor-space-isomorphic-to-Hilbert-cube:

Cantor-space measurable-isomorphic Hilbert-cube

proof -

Isomorphism between \mathcal{C} and $[0, 1]$

have Cantor-space-isomorphic-to-01closed: Cantor-space measurable-isomorphic (restrict-space borel $\{0..1::\text{real}\}$)

proof -

have space-Cantor-space: space Cantor-space = ($\prod_E i \in UNIV. \{0,1\}$)

by(simp add: Cantor-space-def space-PiM)

have space-Cantor-space-01[simp]: $0 \leq x_n \leq 1 \wedge x_n \in \{0,1\}$ if $x \in \text{space Cantor-space}$ for x_n

using PiE-mem[OF that[simplified space-Cantor-space],of n]

by auto

have Cantor-minus-abs-cantor: $(\lambda n. |x_n - y_n|) \in \text{space Cantor-space}$ if $x \in \text{space Cantor-space} \wedge y \in \text{space Cantor-space}$ for $x y$

unfolding space-Cantor-space

proof safe

fix n

assume $|x_n - y_n| \neq 0$

then consider $x_n = 0 \wedge y_n = 1 \mid x_n = 1 \wedge y_n = 0$

using space-Cantor-space-01[OF assms(1),of n] space-Cantor-space-01[OF assms(2),of n]

by auto

thus $|x_n - y_n| = 1$

by cases auto

qed simp

define Cantor-to-01 :: $(\text{nat} \Rightarrow \text{real}) \Rightarrow \text{real}$ where

Cantor-to-01 $\equiv (\lambda x. (\sum n. (1/3)^{\wedge}(\text{Suc } n) * x_n))$

Cantor-to-01 is a measurable injective embedding.

have Cantor-to-01-summable'[simp]: summable $(\lambda n. (1/3)^{\wedge}(\text{Suc } n) * x_n)$ if $x \in \text{space Cantor-space}$ for x

proof(rule summable-comparison-test[where $g = \lambda n. (1/3)^{\wedge} n$ and $N=0$])

show norm $((1/3)^{\wedge}(\text{Suc } n) * x_n) \leq (1/3)^{\wedge} n$ for n

```

    using space-Cantor-space-01 [OF that, of n] by auto
  qed simp

  have Cantor-to-01-summable[simp]:  $\bigwedge x. x \in \text{space Cantor-space} \implies \text{summable}$ 
    ( $\lambda n. (1/3)^{\wedge n} * x n$ )
    using Cantor-to-01-summable' by simp

  have Cantor-to-01-subst-summable[simp]:  $\text{summable } (\lambda n. (1/3)^{\wedge n} * (x n - y$ 
n)) if assms:  $x \in \text{space Cantor-space } y \in \text{space Cantor-space}$  for  $x y$ 
  proof (rule summable-comparison-test' [where  $g = \lambda n. (1/3)^{\wedge n}$  and  $N = 0$ ])
    show  $\text{norm } ((1 / 3)^{\wedge n} * (x n - y n)) \leq (1 / 3)^{\wedge n}$  for  $n$ 
      using space-Cantor-space-01 [OF Cantor-minus-abs-cantor [OF assms], of  $n$ ]
      by (auto simp: idom-abs-sgn-class.abs-mult)
  qed simp

  have Cantor-to-01-image:  $\text{Cantor-to-01} \in \text{space Cantor-space} \rightarrow \{0..1\}$ 
  proof
    fix  $x$ 
    assume  $h: x \in \text{space Cantor-space}$ 
    have  $\text{Cantor-to-01 } x \leq (\sum n. (1/3)^{\wedge} (\text{Suc } n))$ 
      unfolding Cantor-to-01-def
      by (rule suminf-le) (use  $h$  Cantor-to-01-summable [OF  $h$ ] in auto)
    also have  $\dots = (\sum n. (1 / 3)^{\wedge} n) - (1::\text{real})$ 
      using suminf-minus-initial-segment [OF complete-algebra-summable-geometric [of
 $1/3::\text{real}$ ], of 1]
      by auto
    finally have  $\text{Cantor-to-01 } x \leq 1$ 
      by (simp add: suminf-geometric [of 1/3])
    moreover have  $0 \leq \text{Cantor-to-01 } x$ 
      unfolding Cantor-to-01-def
      by (rule suminf-nonneg) (use Cantor-to-01-summable [OF  $h$ ]  $h$  in auto)
    ultimately show  $\text{Cantor-to-01 } x \in \{0..1\}$ 
      by simp
  qed

  have Cantor-to-01-measurable:  $\text{Cantor-to-01} \in \text{Cantor-space} \rightarrow_M \text{restrict-space}$ 
     $\text{borel } \{0..1\}$ 
  proof (rule measurable-restrict-space2)
    show  $\text{Cantor-to-01} \in \text{borel-measurable Cantor-space}$ 
      unfolding Cantor-to-01-def
    proof (rule borel-measurable-suminf)
      fix  $n$ 
      have  $(\lambda x. x n) \in \text{Cantor-space} \rightarrow_M \text{restrict-space borel } \{0, 1\}$ 
        by (simp add: Cantor-space-def)
      hence  $(\lambda x. x n) \in \text{borel-measurable Cantor-space}$ 
        by (simp add: measurable-restrict-space2-iff)
      thus  $(\lambda x. (1 / 3)^{\wedge} \text{Suc } n * x n) \in \text{borel-measurable Cantor-space}$ 
        by simp
    qed
  qed
  qed (rule Cantor-to-01-image)

```

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have Cantor-to-01-inj: inj-on Cantor-to-01 (space Cantor-space)
and Cantor-to-01-preserves-sets: A ∈ sets Cantor-space ⇒ Cantor-to-01 ‘ A
∈ sets (restrict-space borel {0..1}) for A
proof –
  have sets-Cantor: sets Cantor-space = sets (borel-of (product-topology (λ-.
subtopology euclidean {0,1}) UNIV))
    (is ?lhs = -)
  proof –
    have ?lhs = sets (ΠM i ∈ UNIV. borel-of (subtopology euclidean {0,1}))
      by (simp add: Cantor-space-def borel-of-euclidean borel-of-subtopology)
    thus ?thesis
      by(auto intro!: sets-PiM-equal-borel-of-second-countable-subtopology Polish-space-imp-second-countable[of euclideanreal])
  qed
  have s:space Cantor-space = topspace (product-topology (λ-. subtopology euclidean {0,1}) UNIV)
    by(simp add: space-Cantor-space)

let ?d = λx y::real. if (x = 0 ∨ x = 1) ∧ (y = 0 ∨ y = 1) then dist x y else
0
interpret d01: Metric-space {0,1::real} ?d
  by(auto simp: Metric-space-def)
have d01: d01.mtopology = top-of-set {0,1} d01.mcomplete
proof –
  interpret Metric-space {0,1} dist
    by (simp add: Met-TC.subspace)
  have d01.mtopology = mtopology
    by(auto intro!: Metric-space-eq-mtopology simp: Metric-space-def metric-space-class.dist-commute)
  also have ... = top-of-set {0,1}
    by(auto intro!: Submetric.mtopology-submetric[of UNIV dist {0,1::real},simplified]
simp: Submetric-def Metric-space-def Submetric-axioms-def dist-real-def)
  finally show d01.mtopology = top-of-set {0,1} .
  show d01.mcomplete
    using Metric-space-eq-mcomplete[OF d01.Metric-space-axioms,of dist]
d01.compact-space-eq-Bolzano-Weierstrass d01.compact-space-imp-mcomplete finite.emptyI
finite-subset by blast
  qed
interpret pd: Product-metric 1/3 UNIV id id λ-. {0,1::real} λ-. ?d 1
  by(auto intro!: product-metric-natI d01.Metric-space-axioms)
have mpd-top: pd.Product-metric.mtopology = Cantor-space-topology
  by(auto simp: pd.Product-metric-mtopology-eq[symmetric] d01 intro!: product-topology-cong)
have pd-mcomplete: pd.Product-metric.mcomplete
  by(auto intro!: pd.mcomplete-Mi-mcomplete-M d01)
interpret m01: Submetric UNIV dist {0..1::real}
by(simp add: Submetric-def Submetric-axioms-def Met-TC.Metric-space-axioms)
have restrict-space borel {0..1} = borel-of m01.sub.mtopology

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    by (simp add: borel-of-euclidean borel-of-subtopology m01.mtopology-submetric)
    have pd-def: pd.product-dist x y = ( $\sum n. (1/3)^{\wedge n} * |x n - y n|$ ) if x ∈ space
    Cantor-space y ∈ space Cantor-space for x y
    using space-Cantor-space-01[OF that(1)] space-Cantor-space-01[OF that(2)]
    that by(auto simp: product-dist-def dist-real-def)
    have 1: |Cantor-to-01 x - Cantor-to-01 y| ≤ pd.product-dist x y (is ?lhs ≤
    ?rhs) if x ∈ space Cantor-space y ∈ space Cantor-space for x y
    proof -
      have ?lhs = |( $\sum n. (1/3)^{\wedge(Suc n)} * x n - (1/3)^{\wedge(Suc n)} * y n$ )|
      using that by(simp add: suminf-diff Cantor-to-01-def)
      also have ... = | $\sum n. (1/3)^{\wedge(Suc n)} * (x n - y n)$ |
      by (simp add: right-diff-distrib)
      also have ... ≤ ( $\sum n. |(1/3)^{\wedge(Suc n)} * (x n - y n)|$ )
      proof(rule summable-rabs)
        have ( $\lambda n. |(1/3)^{\wedge(Suc n)} * (x n - y n)|$ ) = ( $\lambda n. (1/3)^{\wedge(Suc n)} * |x
        n - y n|$ )
        by (simp add: abs-mult-pos mult.commute)
        moreover have summable ...
        using Cantor-minus-abs-cantor[OF that] by simp
        ultimately show summable ( $\lambda n. |(1/3)^{\wedge(Suc n)} * (x n - y n)|$ ) by
    simp
      qed
      also have ... = ( $\sum n. (1/3)^{\wedge(Suc n)} * |x n - y n|$ )
      by (simp add: abs-mult-pos mult.commute)
      also have ... ≤ pd.product-dist x y
      unfolding pd-def[OF that]
      by(rule suminf-le) (use Cantor-minus-abs-cantor[OF that] in auto)
      finally show ?thesis .
    qed

    have 2: |Cantor-to-01 x - Cantor-to-01 y| ≥ 1 / 9 * pd.product-dist x y (is
    ?lhs ≤ ?rhs) if x ∈ space Cantor-space y ∈ space Cantor-space for x y
    proof(cases x = y)
      case True
      then show ?thesis
      using pd.Product-metric.zero[of x y] that by(simp add: space-Cantor-space)
    next
      case False
      then obtain n' where x n' ≠ y n' by auto
      define n where n ≡ Min {n. n ≤ n' ∧ x n ≠ y n}
      have n ≤ n'
      using ⟨x n' ≠ y n'⟩ n-def by fastforce
      have x n ≠ y n
      using ⟨x n' ≠ y n'⟩ linorder-class.Min-in[of {n. n ≤ n' ∧ x n ≠ y n}]
      by(auto simp: n-def)
      have  $\forall i < n. x i = y i$ 
      proof safe
        fix i
        assume i < n

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show  $x\ i = y\ i$ 
proof(rule ccontr)
  assume  $x\ i \neq y\ i$ 
  then have  $i \in \{n. n \leq n' \wedge x\ n \neq y\ n\}$ 
    using  $\langle n \leq n' \rangle \langle i < n \rangle$  by auto
  thus False
    using  $\langle i < n \rangle$  linorder-class.Min-gr-iff[of  $\{n. n \leq n' \wedge x\ n \neq y\ n\}$   $i$ ]
 $\langle x\ n' \neq y\ n' \rangle$ 
    by(auto simp: n-def)
  qed
qed
have  $u1: (1/3) \wedge (Suc\ n) * (1/2) \leq |Cantor-to-01\ x - Cantor-to-01\ y|$ 
proof -
  have  $(1/3) \wedge (Suc\ n) * (1/2) \leq |(\sum m. (1/3) \wedge (Suc\ (m + Suc\ n)) * (x\ (m + Suc\ n) - y\ (m + Suc\ n))) + (1/3) \wedge Suc\ n * (x\ n - y\ n)|$ 
proof -
  have  $(1/3) \wedge Suc\ n - (1/3) \wedge (n + 2) * 3/2 \leq (1/3) \wedge Suc\ n - |(\sum m. (1/3) \wedge Suc\ (m + Suc\ n) * (y\ (m + Suc\ n) - x\ (m + Suc\ n)))|$ 
proof -
  have  $|(\sum m. (1/3) \wedge Suc\ (m + Suc\ n) * (y\ (m + Suc\ n) - x\ (m + Suc\ n)))| \leq (1/3) \wedge (n + 2) * 3/2$ 
    (is ?lhs  $\leq$  -)
proof -
  have ?lhs  $\leq (\sum m. |(1/3) \wedge Suc\ (m + Suc\ n) * (y\ (m + Suc\ n) - x\ (m + Suc\ n))|)$ 
    apply(rule summable-rabs,rule summable-ignore-initial-segment[of - Suc\ n])
  using Cantor-minus-abs-cantor[OF that(2,1)] by(simp add: abs-mult)
  also have  $\dots = (\sum m. (1/3) \wedge Suc\ (m + Suc\ n) * |y\ (m + Suc\ n) - x\ (m + Suc\ n)|)$ 
    by(simp add: abs-mult)
  also have  $\dots \leq (\sum m. (1/3) \wedge Suc\ (m + Suc\ n))$ 
    apply(rule suminf-le)
  using space-Cantor-space-01[OF Cantor-minus-abs-cantor[OF that(2,1)]]
    apply simp
    apply(rule summable-ignore-initial-segment[of - Suc\ n])
  using Cantor-minus-abs-cantor[OF that(2,1)] by auto
  also have  $\dots = (\sum m. (1/3) \wedge (m + Suc\ (Suc\ n)) * 1)$  by simp
  also have  $\dots = (1/3) \wedge (n + 2) * 3/(2::real)$ 
    by(simp only: pd.nsum-of-rK[of Suc\ (Suc\ n)],simp)
  finally show ?thesis .
qed
thus ?thesis by simp
qed
also have  $\dots = |(1/3) \wedge Suc\ n * (x\ n - y\ n)| - |(\sum m. (1/3) \wedge Suc\ (m + Suc\ n) * (y\ (m + Suc\ n) - x\ (m + Suc\ n)))|$ 
    using  $\langle x\ n \neq y\ n \rangle$  space-Cantor-space-01[OF Cantor-minus-abs-cantor[OF that],of n] by(simp add: abs-mult)

```

also have ... $\leq |(1 / 3) \wedge \text{Suc } n * (x \ n - y \ n) - (\sum m. (1 / 3) \wedge \text{Suc } (m + \text{Suc } n) * (y (m + \text{Suc } n) - x (m + \text{Suc } n)))|$
by *simp*
also have ... $= |(1 / 3) \wedge \text{Suc } n * (x \ n - y \ n) + (\sum m. (1 / 3) \wedge \text{Suc } (m + \text{Suc } n) * (x (m + \text{Suc } n) - y (m + \text{Suc } n)))|$
proof -
have $(\sum m. (1 / 3) \wedge \text{Suc } (m + \text{Suc } n) * (x (m + \text{Suc } n) - y (m + \text{Suc } n))) = (\sum m. - ((1 / 3) \wedge \text{Suc } (m + \text{Suc } n) * (y (m + \text{Suc } n) - x (m + \text{Suc } n))))$
proof -
{ **fix** *nn* :: *nat*
have $\bigwedge r \ ra \ rb. - ((- (r::\text{real}) + ra) / (1 / rb)) = (- ra + r) / (1 / rb)$
by (*simp add: left-diff-distrib*)
then have $- ((y (\text{Suc } (n + nn)) + - x (\text{Suc } (n + nn))) * (1 / 3) \wedge \text{Suc } (\text{Suc } (n + nn))) = (x (\text{Suc } (n + nn)) + - y (\text{Suc } (n + nn))) * (1 / 3) \wedge \text{Suc } (\text{Suc } (n + nn))$
by *fastforce*
then have $- ((1 / 3) \wedge \text{Suc } (nn + \text{Suc } n) * (y (nn + \text{Suc } n) - x (nn + \text{Suc } n))) = (1 / 3) \wedge \text{Suc } (nn + \text{Suc } n) * (x (nn + \text{Suc } n) - y (nn + \text{Suc } n))$
by (*simp add: add.commute mult.commute*) }
then show ?thesis
by *presburger*
qed
also have ... $= - (\sum m. (1 / 3) \wedge \text{Suc } (m + \text{Suc } n) * (y (m + \text{Suc } n) - x (m + \text{Suc } n)))$
apply(*rule suminf-minus*)
apply(*rule summable-ignore-initial-segment[of - Suc n]*)
using *that* **by** *simp*
finally show ?thesis **by** *simp*
qed
also have ... $= |(\sum m. (1 / 3) \wedge \text{Suc } (m + \text{Suc } n) * (x (m + \text{Suc } n) - y (m + \text{Suc } n))) + (1 / 3) \wedge \text{Suc } n * (x \ n - y \ n)|$
using *1* **by** *simp*
finally show ?thesis **by** *simp*
qed
also have ... $= |(\sum m. (1/3) \wedge (\text{Suc } (m + \text{Suc } n)) * (x (m + \text{Suc } n) - y (m + \text{Suc } n))) + (\sum m < \text{Suc } n. (1/3) \wedge (\text{Suc } m) * (x \ m - y \ m))|$
using $\langle \forall i < n. x \ i = y \ i \rangle$ **by** *auto*
also have ... $= |\sum n. (1/3) \wedge (\text{Suc } n) * (x \ n - y \ n)|$
proof -
have $(\sum n. (1 / 3) \wedge \text{Suc } n * (x \ n - y \ n)) = (\sum m. (1 / 3) \wedge \text{Suc } (m + \text{Suc } n) * (x (m + \text{Suc } n) - y (m + \text{Suc } n))) + (\sum m < \text{Suc } n. (1 / 3) \wedge \text{Suc } m * (x \ m - y \ m))$
by(*rule suminf-split-initial-segment*) (*use that in simp*)
thus ?thesis **by** *simp*
qed
also have ... $= |(\sum n. (1/3) \wedge (\text{Suc } n) * x \ n - (1/3) \wedge (\text{Suc } n) * y \ n)|$

```

    by (simp add: right-diff-distrib)
  also have ... = |Cantor-to-01 x - Cantor-to-01 y|
    using that by (simp add: suminf-diff Cantor-to-01-def)
  finally show ?thesis .
qed
have u2: (1/9) * pd.product-dist x y ≤ (1/3) ^ (Suc n) * (1/2)
proof -
  have pd.product-dist x y = (∑ m. (1/3) ^ m * |x m - y m|)
  by (simp add: that pd-def)
  also have ... = (∑ m. (1/3) ^ (m + n) * |x (m + n) - y (m + n)|) +
(∑ m < n. (1/3) ^ m * |x m - y m|)
  using Cantor-minus-abs-cantor[OF that] by (auto intro!: suminf-split-initial-segment)
  also have ... = (∑ m. (1/3) ^ (m + n) * |x (m + n) - y (m + n)|)
  using ⟨∀ i < n. x i = y i⟩ by simp
  also have ... ≤ (∑ m. (1/3) ^ (m + n))
  using space-Cantor-space-01[OF Cantor-minus-abs-cantor[OF that]]
Cantor-minus-abs-cantor[OF that]
  by (auto intro!: suminf-le summable-ignore-initial-segment[of - n])
  also have ... = (1 / 3) ^ n * (3 / 2)
  using pd.nsum-of-rK[of n] by auto
  finally show ?thesis
  by auto
qed
from u1 u2 show ?thesis by simp
qed

have inj: inj-on Cantor-to-01 (space Cantor-space)
proof
  fix x y
  assume h: x ∈ space Cantor-space y ∈ space Cantor-space
  Cantor-to-01 x = Cantor-to-01 y
  then have pd.product-dist x y = 0
  using 2[OF h(1,2)] pd.Product-metric.nonneg[of x y]
  by simp
  thus x = y
  using pd.Product-metric.zero[of x y] h(1,2)
  by (simp add: space-Cantor-space)
qed

have closed: closedin m01.sub.mtopology (Cantor-to-01 ` (space Cantor-space))
  unfolding m01.sub.metric-closedin-iff-sequentially-closed
proof safe
  show a ∈ space Cantor-space ⇒ Cantor-to-01 a ∈ {0..1} for a
  using Cantor-to-01-image by auto
next
  fix xn x
  assume h: range xn ⊆ Cantor-to-01 ` space Cantor-space limitin m01.sub.mtopology
  xn x sequentially
  have ∧n. xn n ∈ {0..1}

```

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    using h(1) measurable-space[OF Cantor-to-01-measurable]
    by (metis (no-types, lifting) UNIV-I atLeastAtMost-borel image-subset-iff
space-restrict-space2 subsetD)
  with h(2) have xnC:m01.sub.MCauchy xn
    by(auto intro!: m01.sub.convergent-imp-MCauchy)
  have  $\forall n. \exists x \in \text{space Cantor-space}. xn\ n = \text{Cantor-to-01 } x$  using h(1) by
auto
  then obtain yn where hx: $\bigwedge n. yn\ n \in \text{space Cantor-space} \wedge n. xn\ n =$ 
Cantor-to-01 (yn n) by metis
  have pd.Product-metric.MCauchy yn
    unfolding pd.Product-metric.MCauchy-def
  proof safe
    fix  $\varepsilon$ 
    assume (0 :: real) <  $\varepsilon$ 
    hence 0 <  $\varepsilon / 9$  by auto
    then obtain N' where  $\bigwedge n\ m. n \geq N' \implies m \geq N' \implies |xn\ n - xn\ m|$ 
<  $\varepsilon / 9$ 
      using xnC m01.sub.MCauchy-def xnC unfolding dist-real-def by blast
    thus  $\exists N. \forall n\ n'. N \leq n \longrightarrow N \leq n' \longrightarrow pd.\text{product-dist } (yn\ n) (yn\ n')$ 
<  $\varepsilon$ 
      using order.strict-trans1[OF 2[OF hx(1) hx(1)],of - -  $\varepsilon/9$ ] hx(1)
      by(auto intro!: exI[where x=N'] simp: hx(2) space-Cantor-space)
    qed(use hx space-Cantor-space in auto)
  then obtain y where y:limitin pd.Product-metric.mtopology yn y sequentially
    using pd-mcomplete pd.Product-metric.mcomplete-def by blast
  hence y  $\in \text{space Cantor-space}$ 
    by (simp add: pd.Product-metric.limitin-mspace space-Cantor-space)
  have limitin m01.sub.mtopology xn (Cantor-to-01 y) sequentially
    unfolding m01.sub.limit-metric-sequentially
  proof safe
    show Cantor-to-01 y  $\in \{0..1\}$ 
    using h(1) funcset-image[OF Cantor-to-01-image]  $\langle y \in \text{space Cantor-space} \rangle$ 
  by blast
  next
    fix  $\varepsilon$ 
    assume (0 :: real) <  $\varepsilon$ 
    then obtain N where  $\bigwedge n. n \geq N \implies pd.\text{product-dist } (yn\ n) y < \varepsilon \wedge n.$ 
n  $\geq N \implies yn\ n \in UNIV \rightarrow_E \{0, 1\}$ 
      using y by(fastforce simp: pd.Product-metric.limit-metric-sequentially)
    with  $\langle \bigwedge n. xn\ n \in \{0..1\} \rangle$  show  $\exists N. \forall n \geq N. xn\ n \in \{0..1\} \wedge dist (xn$ 
n) (Cantor-to-01 y) <  $\varepsilon$ 
      by(auto intro!: exI[where x=N] order.strict-trans1[OF 1[OF hx(1)  $\langle y$ 
 $\in \text{space Cantor-space} \rangle$ ]] simp: submetric-def  $\langle 0 < \varepsilon \rangle$  hx(2) dist-real-def)
    qed
    hence Cantor-to-01 y = x
      using h(2) by(auto dest: m01.sub.limitin-metric-unique)
  with  $\langle y \in \text{space Cantor-space} \rangle$  show x  $\in \text{Cantor-to-01 'space Cantor-space}$ 
    by auto
  qed

```



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have open-map:open-map pd.Product-metric.mtopology (subtopology m01.sub.mtopology
(Cantor-to-01 ' (space Cantor-space))) Cantor-to-01
proof –
have open-map (mtopology-of pd.Product-metric.Self) (subtopology (mtopology-of
m01.sub.Self) (Cantor-to-01 ' mspace pd.Product-metric.Self)) Cantor-to-01
proof(rule Metric-space-open-map-from-dist)
fix x ε
assume (0 :: real) < ε x ∈ mspace pd.Product-metric.Self
then have x ∈ (UNIV :: nat set) →E {0, 1::real}
by simp
show ∃δ>0. ∀y∈mspace pd.Product-metric.Self. mdist m01.sub.Self
(Cantor-to-01 x) (Cantor-to-01 y) < δ → mdist pd.Product-metric.Self x y < ε
unfolding pd.Product-metric.mspace-Self pd.Product-metric.mdist-Self
m01.sub.mdist-Self
proof(safe intro!: exI[where x=ε/9])
fix y
assume h:y ∈ (UNIV :: nat set) →E {0, 1::real} dist (Cantor-to-01 x)
(Cantor-to-01 y) < ε / 9
then have sc:x ∈ space Cantor-space y ∈ space Cantor-space
using ⟨x ∈ UNIV →E {0, 1}⟩ by(simp-all add: space-Cantor-space)
have |Cantor-to-01 x – Cantor-to-01 y| < ε / 9
using h by(simp add: dist-real-def)
with 2[OF sc] show pd.product-dist x y < ε
by simp
qed (use ⟨ε > 0⟩ in auto)
qed(use Cantor-to-01-image space-Cantor-space in auto)
thus ?thesis
by (simp add: mtopology-of-def space-Cantor-space)
qed
have Cantor-to-01 ' A ∈ sets (restrict-space borel {0..1}) if A ∈ sets Can-
tor-space for A
using open-map-preserves-sets'[of pd.Product-metric.mtopology m01.sub.mtopology
Cantor-to-01 A] borel-of-closed[OF closed] inj sets-Cantor open-map that mpd-top
⟨restrict-space borel {0..1} = borel-of m01.sub.mtopology⟩
by (simp add: space-Cantor-space)
with inj show inj-on Cantor-to-01 (space Cantor-space)
andA ∈ sets Cantor-space ⇒ Cantor-to-01 ' A ∈ sets (restrict-space borel
{0..1})
by simp-all
qed

```

Next, we construct measurable embedding from $[0, 1]$ to $0, 1^{\mathbb{N}}$.

```

define to-Cantor-from-01 :: real ⇒ nat ⇒ real where
to-Cantor-from-01 ≡ (λr n. if r = 1 then 1 else real-of-int (⌊2∧(Suc n) * r⌋
mod 2))

```

to-Cantor-from-01 is a measurable injective embedding into Cantor space.

```

have to-Cantor-from-01-image': to-Cantor-from-01 r n ∈ {0,1} for r n

```

unfolding *to-Cantor-from-01-def* **by** *auto*
have *to-Cantor-from-01-image'*: $\bigwedge r n. 0 \leq \text{to-Cantor-from-01 } r n \bigwedge r n.$
to-Cantor-from-01 } r n \leq 1
by (*auto simp add: to-Cantor-from-01-def*)
have *to-Cantor-from-01-image*: $\text{to-Cantor-from-01} \in \{0..1\} \rightarrow \text{space Cantor-space}$
using *to-Cantor-from-01-image'* **by**(*auto simp: space-Cantor-space*)
have *to-Cantor-from-01-measurable*:
to-Cantor-from-01} \in \text{restrict-space borel } \{0..1\} \rightarrow_M \text{Cantor-space}
unfolding *to-Cantor-from-01-def Cantor-space-def*
by(*auto intro!: measurable-restrict-space3 measurable-abs-UNIV*)
have *to-Cantor-from-01-summable[simp]*:
summable ($\lambda n. (1/2)^{\wedge} n * \text{to-Cantor-from-01 } r n$) **for** *r*
proof(*rule summable-comparison-test'*[**where** $g = \lambda n. (1/2)^{\wedge} n$])
show $\text{norm } ((1/2)^{\wedge} n * \text{to-Cantor-from-01 } r n) \leq (1/2)^{\wedge} n$ **for** *n*
using *to-Cantor-from-01-image'*[*of r n*] **by** *auto*
qed *simp*

have *to-Cantor-from-sumn'*: $(\sum i < n. (1/2)^{\wedge} (\text{Suc } i) * \text{to-Cantor-from-01 } r i) \leq$
r
 $r - (\sum i < n. (1/2)^{\wedge} (\text{Suc } i) * \text{to-Cantor-from-01 } r i) < (1/2)^{\wedge} n$
 $\text{to-Cantor-from-01 } r n = 1 \iff (1/2)^{\wedge} (\text{Suc } n) \leq r - (\sum i < n. (1/2)^{\wedge} (\text{Suc } i) * \text{to-Cantor-from-01 } r i)$
 $\text{to-Cantor-from-01 } r n = 0 \iff r - (\sum i < n. (1/2)^{\wedge} (\text{Suc } i) * \text{to-Cantor-from-01 } r i) < (1/2)^{\wedge} (\text{Suc } n)$ **if** *assms: r \in \{0..<1\}* **for** *r n*
proof –
let *?f = to-Cantor-from-01 r*
have *f-simp*: $?f l = \text{real-of-int } (\lfloor 2^{\wedge} (\text{Suc } l) * r \rfloor \text{ mod } 2)$ **for** *l*
using *assms* **by**(*simp add: to-Cantor-from-01-def*)
define *S* **where** $S = (\lambda n. \sum i < n. (1/2)^{\wedge} (\text{Suc } i) * ?f i)$
have *SSuc*: $S (\text{Suc } k) = S k + (1/2)^{\wedge} (\text{Suc } k) * \text{to-Cantor-from-01 } r k$ **for** *k*
by(*simp add: S-def*)
have *Sfloor*: $\lfloor 2^{\wedge} (\text{Suc } m) * (l - S m) \rfloor \text{ mod } 2 = \lfloor 2^{\wedge} (\text{Suc } m) * l \rfloor \text{ mod } 2$ **for** *l*
m
proof –
have $\exists z. 2^{\wedge} (\text{Suc } m) * ((1/2)^{\wedge} (\text{Suc } k) * ?f k) = 2 * \text{real-of-int } z$ **if** $k < m$
for *k*
proof –
have $0: (2::\text{real})^{\wedge} m * (1/2)^{\wedge} k = 2 * 2^{\wedge} (m-k-1)$
using *that* **by** (*simp add: power-diff-conv-inverse*)
consider $?f k = 0 \mid ?f k = 1$
using *to-Cantor-from-01-image'*[*of r k*] **by** *auto*
thus *?thesis*
apply *cases* **using** *that 0* **by** *auto*
qed
then **obtain** *z* **where** $\bigwedge k. k < m \implies 2^{\wedge} (\text{Suc } m) * ((1/2)^{\wedge} (\text{Suc } k) * ?f k) = 2 * \text{real-of-int } (z k)$
by *metis*
hence $S m: 2^{\wedge} (\text{Suc } m) * S m = \text{real-of-int } (2 * (\sum k < m. (z k)))$

```

    by(auto simp: S-def sum-distrib-left)
    have  $\lfloor 2^{\wedge}(Suc\ m) * (l - S\ m) \rfloor \bmod 2 = \lfloor 2^{\wedge}(Suc\ m) * l - 2^{\wedge}(Suc\ m) * S\ m \rfloor \bmod 2$ 
      by (simp add: right-diff-distrib)
    also have ... =  $\lfloor 2^{\wedge}(Suc\ m) * l \rfloor \bmod 2$ 
      unfolding Sm
      by(simp only: floor-diff-of-int) presburger
    finally show ?thesis .
  qed
  have  $S\ n \leq r \wedge r - S\ n < (1/2)^{\wedge}n \wedge (?f\ n = 1 \iff (1/2)^{\wedge}(Suc\ n) \leq r - S\ n) \wedge (?f\ n = 0 \iff r - S\ n < (1/2)^{\wedge}(Suc\ n))$ 
  proof(induction n)
    case 0
    then show ?case
      using assms by(auto simp: S-def to-Cantor-from-01-def) linarith+
  next
    case (Suc n)
    hence ih:  $S\ n \leq r \wedge r - S\ n < (1/2)^{\wedge}n$ 
      ?f n = 1  $\implies (1/2)^{\wedge}Suc\ n \leq r - S\ n$ 
      ?f n = 0  $\implies r - S\ n < (1/2)^{\wedge}Suc\ n$ 
      by simp-all
    have SSuc':  $?f\ n = 0 \wedge S\ (Suc\ n) = S\ n \vee ?f\ n = 1 \wedge S\ (Suc\ n) = S\ n + (1/2)^{\wedge}(Suc\ n)$ 
      using to-Cantor-from-01-image'[of r n] by(simp add: SSuc)
    have goal1:  $S\ (Suc\ n) \leq r$ 
      using SSuc' ih(1) ih(3) by auto
    have goal2:  $r - S\ (Suc\ n) < (1/2)^{\wedge}Suc\ n$ 
      using SSuc' ih(4) ih(2) by auto
    have goal3-1:  $(1/2)^{\wedge}Suc\ (Suc\ n) \leq r - S\ (Suc\ n)$  if ?f (Suc n) = 1
    proof(rule ccontr)
      assume  $\neg (1/2)^{\wedge}Suc\ (Suc\ n) \leq r - S\ (Suc\ n)$ 
      then have  $r - S\ (Suc\ n) < (1/2)^{\wedge}Suc\ (Suc\ n)$  by simp
      hence h:  $2^{\wedge}Suc\ (Suc\ n) * (r - S\ (Suc\ n)) < 1$ 
        using mult-less-cancel-left-pos[of  $2^{\wedge}Suc\ (Suc\ n)$   $r - S\ (Suc\ n)$   $(1/2)^{\wedge}Suc\ (Suc\ n)$ ]
        by (simp add: power-one-over)
      moreover have  $0 \leq 2^{\wedge}Suc\ (Suc\ n) * (r - S\ (Suc\ n))$ 
        using goal1 by simp
      ultimately have  $\lfloor 2^{\wedge}Suc\ (Suc\ n) * (r - S\ (Suc\ n)) \rfloor = 0$ 
        by linarith
      thus False
        using that[simplified f-simp] Sfloor[of Suc n r]
        by fastforce
    qed
    have goal3-2:  $?f\ (Suc\ n) = 1$  if  $(1/2)^{\wedge}Suc\ (Suc\ n) \leq r - S\ (Suc\ n)$ 
    proof -
      have  $1 \leq 2^{\wedge}Suc\ (Suc\ n) * (r - S\ (Suc\ n))$ 
        using that[simplified f-simp] mult-le-cancel-left-pos[of  $2^{\wedge}Suc\ (Suc\ n)$   $(1/2)^{\wedge}Suc\ (Suc\ n)$   $r - S\ (Suc\ n)$ ]

```

```

    by (simp add: power-one-over)
    moreover have  $2^{\wedge} \text{Suc } (\text{Suc } n) * (r - S (\text{Suc } n)) < 2$ 
    using mult-less-cancel-left-pos[of  $2^{\wedge} \text{Suc } (\text{Suc } n)$   $r - S (\text{Suc } n)$   $(1 / 2)$ 
 $^{\wedge} \text{Suc } n$ ] goal2
    by (simp add: power-one-over)
    ultimately have  $[2^{\wedge} \text{Suc } (\text{Suc } n) * (r - S (\text{Suc } n))] = 1$ 
    by linarith
    thus ?thesis
    using Sfloor[of  $\text{Suc } n$   $r$ ] by(auto simp: f-simp)
qed
have goal4-1:  $r - S (\text{Suc } n) < (1 / 2)^{\wedge} \text{Suc } (\text{Suc } n)$  if  $?f (\text{Suc } n) = 0$ 
proof(rule ccontr)
  assume  $\neg r - S (\text{Suc } n) < (1 / 2)^{\wedge} \text{Suc } (\text{Suc } n)$ 
  then have  $(1 / 2)^{\wedge} \text{Suc } (\text{Suc } n) \leq r - S (\text{Suc } n)$  by simp
  hence  $1 \leq 2^{\wedge} \text{Suc } (\text{Suc } n) * (r - S (\text{Suc } n))$ 
  using mult-le-cancel-left-pos[of  $2^{\wedge} \text{Suc } (\text{Suc } n)$   $(1 / 2)^{\wedge} \text{Suc } (\text{Suc } n)$   $r$ 
 $- S (\text{Suc } n)$ ]
  by (simp add: power-one-over)
  moreover have  $2^{\wedge} \text{Suc } (\text{Suc } n) * (r - S (\text{Suc } n)) < 2$ 
  using mult-less-cancel-left-pos[of  $2^{\wedge} \text{Suc } (\text{Suc } n)$   $r - S (\text{Suc } n)$   $(1 / 2)$ 
 $^{\wedge} \text{Suc } n$ ] goal2
  by (simp add: power-one-over)
  ultimately have  $[2^{\wedge} \text{Suc } (\text{Suc } n) * (r - S (\text{Suc } n))] = 1$ 
  by linarith
  thus False
  using that Sfloor[of  $\text{Suc } n$   $r$ ] by(auto simp: f-simp)
qed
have goal4-2:  $?f (\text{Suc } n) = 0$  if  $r - S (\text{Suc } n) < (1 / 2)^{\wedge} \text{Suc } (\text{Suc } n)$ 
proof -
  have  $h: 2^{\wedge} \text{Suc } (\text{Suc } n) * (r - S (\text{Suc } n)) < 1$ 
  using mult-less-cancel-left-pos[of  $2^{\wedge} \text{Suc } (\text{Suc } n)$   $r - S (\text{Suc } n)$   $(1 / 2)$ 
 $^{\wedge} \text{Suc } (\text{Suc } n)$ ] that
  by (simp add: power-one-over)
  moreover have  $0 \leq 2^{\wedge} \text{Suc } (\text{Suc } n) * (r - S (\text{Suc } n))$ 
  using goal1 by simp
  ultimately have  $[2^{\wedge} \text{Suc } (\text{Suc } n) * (r - S (\text{Suc } n))] = 0$ 
  by linarith
  thus ?thesis
  using Sfloor[of  $\text{Suc } n$   $r$ ] by(auto simp: f-simp)
qed
show ?case
  using goal1 goal2 goal3-1 goal3-2 goal4-1 goal4-2 by blast
qed
thus  $(\sum_{i < n}. (1/2)^{\wedge} (\text{Suc } i) * \text{to-Cantor-from-01 } r i) \leq r$ 
and  $r - (\sum_{i < n}. (1/2)^{\wedge} (\text{Suc } i) * \text{to-Cantor-from-01 } r i) < (1/2)^{\wedge} n$ 
and  $\text{to-Cantor-from-01 } r n = 1 \iff (1/2)^{\wedge} (\text{Suc } n) \leq r - (\sum_{i < n}. (1/2)^{\wedge} (\text{Suc } i) * \text{to-Cantor-from-01 } r i)$ 
and  $\text{to-Cantor-from-01 } r n = 0 \iff r - (\sum_{i < n}. (1/2)^{\wedge} (\text{Suc } i) * \text{to-Cantor-from-01 } r i) < (1/2)^{\wedge} (\text{Suc } n)$ 

```

by(*simp-all add: S-def*)
qed
have *to-Cantor-from-sumn*: $(\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i) \leq r$
r
 $r - (\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i) \leq (1/2) \wedge n$
 $to-Cantor-from-01\ r\ n = 1 \iff (1/2) \wedge (Suc\ n) \leq r - (\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i)$
 $to-Cantor-from-01\ r\ n = 0 \iff r - (\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i) < (1/2) \wedge (Suc\ n)$ **if** *assms*: $r \in \{0..1\}$ **for** *r n*
proof –
have *nsum*: $(\sum i < n. (1/2) \wedge (Suc\ i)) = 1 - (1 / (2::real)) \wedge n$
using *one-diff-power-eq*[*of 1 / (2::real) n*] **by**(*auto simp: sum-divide-distrib[symmetric]*)

consider $r = 1 \mid r \in \{0..<1\}$ **using** *assms* **by** *fastforce*
hence $(\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i) \leq r \wedge r - (\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i) \leq (1/2) \wedge n \wedge (to-Cantor-from-01\ r\ n = 1 \iff (1/2) \wedge (Suc\ n) \leq r - (\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i)) \wedge (to-Cantor-from-01\ r\ n = 0 \iff r - (\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i) < (1/2) \wedge (Suc\ n))$
proof *cases*
case 1
then show *?thesis*
using *nsum* **by**(*auto simp: to-Cantor-from-01-def*)
next
case 2
from *to-Cantor-from-sumn*'[*OF this*]
show *?thesis*
using *less-eq-real-def* **by** *blast*
qed
thus $(\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i) \leq r$
and $r - (\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i) \leq (1/2) \wedge n$
and $to-Cantor-from-01\ r\ n = 1 \iff (1/2) \wedge (Suc\ n) \leq r - (\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i)$
and $to-Cantor-from-01\ r\ n = 0 \iff r - (\sum i < n. (1/2) \wedge (Suc\ i) * to-Cantor-from-01\ r\ i) < (1/2) \wedge (Suc\ n)$
by *simp-all*
qed

have *to-Cantor-from-sum*: $(\sum n. (1/2) \wedge (Suc\ n) * to-Cantor-from-01\ r\ n) = r$ **if** *assms*: $r \in \{0..1\}$ **for** *r*
proof –
have $1:r \leq (\sum n. (1/2) \wedge (Suc\ n) * to-Cantor-from-01\ r\ n)$
proof –
have $0:r \leq (1 / 2) \wedge n + (\sum n. (1/2) \wedge (Suc\ n) * to-Cantor-from-01\ r\ n)$
for *n*
proof –
have $r \leq (1 / 2) \wedge n + (\sum i < n. (1 / 2) \wedge Suc\ i * to-Cantor-from-01\ r\ i)$
using *to-Cantor-from-sumn*(2)[*OF assms, of n*] **by** *auto*
also have $\dots \leq (1 / 2) \wedge n + (\sum n. (1/2) \wedge (Suc\ n) * to-Cantor-from-01\ r$

n)

```

    using to-Cantor-from-01-image'[of r] by(auto intro!: sum-le-suminf)
    finally show ?thesis .
  qed
  have 00:∃ n0. ∀ n≥n0. (1 / 2) ^ n < r if r>0 for r :: real
  proof -
    obtain n0 where (1 / 2) ^ n0 < r
    using reals-power-lt-ex[of - 2 :: real, OF ‹r>0›] by auto
    thus ?thesis
    using order.strict-trans1[OF power-decreasing[of n0 - 1/2::real]]
    by(auto intro!: exI[where x=n0])
  qed
  show ?thesis
  apply(rule Lim-bounded2[where f=λn. (1 / 2) ^ n + (∑ n. (1/2) ^ (Suc
n)*to-Cantor-from-01 r n) and N=0])
    using 0 00 by(auto simp: LIMSEQ-iff)
  qed
  have 2:(∑ n. (1/2) ^ (Suc n)*to-Cantor-from-01 r n) ≤ r
    using to-Cantor-from-sumn[OF assms] by(auto intro!: suminf-le-const)
  show ?thesis
  using 1 2 by simp
  qed
  have to-Cantor-from-sum': (∑ i<n. (1/2) ^ (Suc i)*to-Cantor-from-01 r i) = r
  - (∑ m. (1/2) ^ (Suc (m + n))*to-Cantor-from-01 r (m + n)) if assms:r ∈ {0..1}
  for r n
    using suminf-minus-initial-segment[of λn. (1 / 2) ^ Suc n * to-Cantor-from-01
r n n] to-Cantor-from-sum[OF assms]
    by auto

  have to-Cantor-from-01-exist0: ∀ n.∃ k≥n. to-Cantor-from-01 r k = 0 if assms:r
  ∈ {0..<1} for r
  proof(rule ccontr)
    assume ¬ (∀ n.∃ k≥n. to-Cantor-from-01 r k = 0)
    then obtain n0 where hn0:
      ∧k. k ≥ n0 ⇒ to-Cantor-from-01 r k = 1
    using to-Cantor-from-01-image'[of r] by auto
    define n where n = Min {i. i ≤ n0 ∧ (∀ k≥i. to-Cantor-from-01 r k = 1)}
    have n0in: n0 ∈ {i. i ≤ n0 ∧ (∀ k≥i. to-Cantor-from-01 r k = 1)}
    using hn0 by auto
    have hn:n ≤ n0 ∧k. k ≥ n ⇒ to-Cantor-from-01 r k = 1
    using n0in Min-in[of {i. i ≤ n0 ∧ (∀ k≥i. to-Cantor-from-01 r k = 1)}]
    by(auto simp: n-def)
    show False
  proof(cases n)
    case 0
    then have r = (∑ n. (1 / 2) ^ Suc n)
    using to-Cantor-from-sum[of r] assms hn(2) by simp
    also have ... = 1
    using nsum-of-r'[of 1/2 1 1] by auto
  end
end

```

```

finally show ?thesis
  using assms by auto
next
  case eqn:(Suc n')
  have to-Cantor-from-01 r n' = 0
  proof(rule ccontr)
    assume to-Cantor-from-01 r n' ≠ 0
    then have to-Cantor-from-01 r n' = 1
      using to-Cantor-from-01-image'[of r n'] by auto
    hence n' ∈ {i. i ≤ n0 ∧ (∀ k ≥ i. to-Cantor-from-01 r k = 1)}
      using hn eqn not-less-eq-eq order-antisym-conv by fastforce
    hence n ≤ n'
      using Min.coboundedI[of {i. i ≤ n0 ∧ (∀ k ≥ i. to-Cantor-from-01 r k =
1)} n']
      by(simp add: n-def)
    thus False
      using eqn by simp
  qed
  hence le1:r - (∑ i < n'. (1 / 2) ^ Suc i * to-Cantor-from-01 r i) < (1 /
2) ^ n
    using to-Cantor-from-sumn'(4)[OF assms, of n'] by (simp add: eqn)
  have r - (∑ i < n'. (1 / 2) ^ Suc i * to-Cantor-from-01 r i) = (1 / 2) ^ n
    (is ?lhs = -)
  proof -
    have ?lhs = (∑ m. (1/2) ^ (m + Suc n') * to-Cantor-from-01 r (m + n'))
      using to-Cantor-from-sum'[of r n'] assms by simp
    also have ... = (∑ m. (1/2) ^ (m + Suc n) * to-Cantor-from-01 r (m + n))
      proof -
        have (∑ n. (1 / 2) ^ (Suc n + Suc n') * to-Cantor-from-01 r (Suc n +
n')) = (∑ m. (1 / 2) ^ (m + Suc n') * to-Cantor-from-01 r (m + n')) - (1 / 2)
^ (0 + Suc n') * to-Cantor-from-01 r (0 + n')
          by(rule suminf-split-head) (auto intro!: summable-ignore-initial-segment)
        thus ?thesis
          using <to-Cantor-from-01 r n' = 0> by(simp add: eqn)
      qed
    also have ... = (∑ m. (1/2) ^ (m + Suc n))
      using hn by simp
    also have ... = (1 / 2) ^ n
      using nsum-of-r'[of 1/2 Suc n 1, simplified] by simp
    finally show ?thesis .
  qed
with le1 show False
  by simp
qed
qed
have to-Cantor-from-01-if-exist0: to-Cantor-from-01 (∑ n. (1 / 2) ^ Suc n *
a n) = a if assms: ∧ n. a n ∈ {0,1} ∀ n. ∃ k ≥ n. a k = 0 for a
proof
  fix n

```

```

have [simp]: summable ( $\lambda n. (1/2)^n * a n$ )
proof(rule summable-comparison-test'[where  $g = \lambda n. (1/2)^n$ ])
  show norm (( $1/2$ )n * a n) ≤ ( $1/2$ )n for n
    using assms(1)[of n] by auto
qed simp
let ?r =  $\sum n. (1/2)^{Suc n} * a n$ 
have ?r ∈ {0..1}
  using assms(1) space-Cantor-space-01[of a, simplified space-Cantor-space]
  nsum-of-r-leq[of 1/2 a 1 1 0]
  by auto
show to-Cantor-from-01 ?r n = a n
proof(rule less-induct)
  fix x
  assume ih:  $y < x \implies to-Cantor-from-01 ?r y = a y$  for y
  have eq1: ?r - ( $\sum i < x. (1/2)^{Suc i} * to-Cantor-from-01 ?r i$ ) = ( $\sum n. (1/2)^{Suc (n+x)} * a (n+x)$ )
    (is ?lhs = ?rhs)
  proof -
    have ?lhs = ( $\sum n. (1/2)^{Suc (n+x)} * a (n+x)$ ) + ( $\sum i < x. (1/2)^{Suc i} * a i$ ) - ( $\sum i < x. (1/2)^{Suc i} * to-Cantor-from-01 ?r i$ )
      using suminf-split-initial-segment[of  $\lambda n. (1/2)^{Suc n} * a n$ ] by simp
    also have ... = ( $\sum n. (1/2)^{Suc (n+x)} * a (n+x)$ ) + ( $\sum i < x. (1/2)^{Suc i} * a i$ ) - ( $\sum i < x. (1/2)^{Suc i} * a i$ )
      using ih by simp
    finally show ?thesis by simp
  qed
define Sn where  $S_n = (\sum n. (1/2)^{Suc (n+x)} * a (n+x))$ 
define Sn' where  $S_n' = (\sum n. (1/2)^{Suc (n+(Suc x))} * a (n+(Suc x)))$ 
have SnSn':  $S_n = (1/2)^{Suc x} * a x + S_n'$ 
  using suminf-split-head[of  $\lambda n. (1/2)^{Suc (n+x)} * a (n+x)$ , OF summable-ignore-initial-segment]
  by(auto simp: Sn-def Sn'-def)
have hsn:  $0 \leq S_n' < (1/2)^{Suc x}$ 
proof -
  show  $0 \leq S_n'$ 
    unfolding Sn'-def
    by(rule suminf-nonneg, rule summable-ignore-initial-segment) (use assms(1) space-Cantor-space-01[of a, simplified space-Cantor-space] in fastforce)+
  next
  have  $\exists n' \geq Suc x. a n' < 1$ 
    using assms by fastforce
  thus  $S_n' < (1/2)^{Suc x}$ 
    using nsum-of-r-le[of 1/2 a 1 Suc x Suc (Suc x)] assms(1) space-Cantor-space-01[of a, simplified space-Cantor-space]
    by(auto simp: Sn'-def)
  qed
have goal1: to-Cantor-from-01 ?r x = 1  $\longleftrightarrow$  a x = 1

```



```

proof –
  have to-Cantor-from-01 ?r x = 1  $\longleftrightarrow$  (1 / 2)  $\wedge$  Suc x  $\leq$  Sn
    using to-Cantor-from-sumn(3)[OF  $\langle$ ?r  $\in$  {0..1} $\rangle$ ] eq1
    by(fastforce simp: Sn-def)
  also have ...  $\longleftrightarrow$  (1 / 2)  $\wedge$  Suc x  $\leq$  (1/2)  $\wedge$  (Suc x) * a x + Sn'
    by(simp add: SnSn')
  also have ...  $\longleftrightarrow$  a x = 1
  proof –
    have a x = 1 if (1 / 2)  $\wedge$  Suc x  $\leq$  (1/2)  $\wedge$  (Suc x) * a x + Sn'
    proof(rule ccontr)
      assume a x  $\neq$  1
      then have a x = 0
        using assms(1) by auto
      hence (1 / 2)  $\wedge$  Suc x  $\leq$  Sn'
        using that by simp
      thus False
        using hsn by auto
    qed
    thus ?thesis
      by(auto simp: hsn)
  qed
finally show ?thesis .
qed
have goal2: to-Cantor-from-01 ?r x = 0  $\longleftrightarrow$  a x = 0
proof –
  have to-Cantor-from-01 ?r x = 0  $\longleftrightarrow$  Sn < (1 / 2)  $\wedge$  Suc x
    using to-Cantor-from-sumn(4)[OF  $\langle$ ?r  $\in$  {0..1} $\rangle$ ] eq1
    by(fastforce simp: Sn-def)
  also have ...  $\longleftrightarrow$  (1/2)  $\wedge$  (Suc x) * a x + Sn' < (1 / 2)  $\wedge$  Suc x
    by(simp add: SnSn')
  also have ...  $\longleftrightarrow$  a x = 0
  proof –
    have a x = 0 if (1/2)  $\wedge$  (Suc x) * a x + Sn' < (1 / 2)  $\wedge$  Suc x
    proof(rule ccontr)
      assume a x  $\neq$  0
      then have a x = 1
        using assms(1) by auto
      thus False
        using that hsn by auto
    qed
    thus ?thesis
      using hsn by auto
  qed
finally show ?thesis .
qed
show to-Cantor-from-01 ?r x = a x
  using goal1 goal2 to-Cantor-from-01-image'[of ?r x] by auto
qed
qed

```

```

have to-Cantor-from-01-sum-of-to-Cantor-from-01: to-Cantor-from-01 ( $\sum n.$ 
( $1 / 2$ )  $\wedge$  Suc  $n$  * to-Cantor-from-01  $r$   $n$ ) = to-Cantor-from-01  $r$  if assms:  $r \in$ 
 $\{0..1\}$  for  $r$ 
proof -
  consider  $r = 1 \mid r \in \{0..<1\}$ 
  using assms by fastforce
  then show ?thesis
  proof cases
  case 1
  then show ?thesis
  using nsum-of-r'[of  $1/2$  1 1]
  by(auto simp: to-Cantor-from-01-def)
  next
  case 2
from to-Cantor-from-01-if-exist0[OF to-Cantor-from-01-image' to-Cantor-from-01-exist0[OF
this]]
  show ?thesis .
qed
qed
have to-Cantor-from-01-inj: inj-on to-Cantor-from-01 (space (restrict-space
borel  $\{0..1\}$ )))
proof
  fix  $x y :: \text{real}$ 
  assume  $x \in \text{space} (\text{restrict-space } \text{borel } \{0..1\})$   $y \in \text{space} (\text{restrict-space } \text{borel}$ 
 $\{0..1\})$ 
  and  $h: \text{to-Cantor-from-01 } x = \text{to-Cantor-from-01 } y$ 
  then have  $xyin: x \in \{0..1\}$   $y \in \{0..1\}$ 
  by simp-all
  show  $x = y$ 
  using to-Cantor-from-sum[OF  $xyin(1)$ ] to-Cantor-from-sum[OF  $xyin(2)$ ]  $h$ 
  by simp
qed
have to-Cantor-from-01-preserves-sets: to-Cantor-from-01 '  $A \in \text{sets } \text{Cantor-space}$ 
if assms:  $A \in \text{sets} (\text{restrict-space } \text{borel } \{0..1\})$  for  $A$ 
proof -
define  $f :: (\text{nat} \Rightarrow \text{real}) \Rightarrow \text{real}$  where  $f \equiv (\lambda x. \sum n. (1/2)^\wedge(\text{Suc } n) * x n)$ 
have  $f\text{-meas}: f \in \text{Cantor-space} \rightarrow_M \text{restrict-space } \text{borel } \{0..1\}$ 
proof -
  have  $f \in \text{borel-measurable } \text{Cantor-space}$ 
  unfolding Cantor-to-01-def f-def
proof(rule borel-measurable-suminf)
  fix  $n$ 
  have  $(\lambda x. x n) \in \text{Cantor-space} \rightarrow_M \text{restrict-space } \text{borel } \{0, 1\}$ 
  by(simp add: Cantor-space-def)
  hence  $(\lambda x. x n) \in \text{borel-measurable } \text{Cantor-space}$ 
  by(simp add: measurable-restrict-space2-iff)
  thus  $(\lambda x. (1 / 2)^\wedge \text{Suc } n * x n) \in \text{borel-measurable } \text{Cantor-space}$ 
  by simp
qed

```

```

moreover have  $0 \leq f x f x \leq 1$  if  $x \in \text{space Cantor-space}$  for  $x$ 
proof –
  have [simp]:summable ( $\lambda n. (1/2)^{\wedge n} * x n$ )
  proof(rule summable-comparison-test'[where  $g = \lambda n. (1/2)^{\wedge n}$ ])
    show norm  $((1 / 2)^{\wedge n} * x n) \leq (1 / 2)^{\wedge n}$  for  $n$ 
    using that by simp
  qed simp
  show  $0 \leq f x$ 
    using that by(auto intro!: suminf-nonneg simp: f-def)
  show  $f x \leq 1$ 
  proof –
    have  $f x \leq (\sum n. (1/2)^{\wedge (Suc n)})$ 
    using that by(auto intro!: suminf-le simp: f-def)
    also have ... = 1
    using nsum-of-r'[of 1/2 1 1] by simp
    finally show ?thesis .
  qed
qed
ultimately show ?thesis
  by(auto intro!: measurable-restrict-space2)
qed
have image-sets:to-Cantor-from-01 ‘ (space (restrict-space borel {0..1}))  $\in$ 
sets Cantor-space
(is ?A  $\in$  -)
proof –
  have ?A  $\subseteq$  space Cantor-space
  using to-Cantor-from-01-image by auto
  have comple-sets:( $\Pi_E i \in UNIV. \{0,1\}$ ) – ?A  $\in$  sets Cantor-space
  proof –
  have eq1: ?A =  $\{\lambda n. 1\} \cup \{x. (\forall n. x n \in \{0,1\}) \wedge (\forall n. \exists k \geq n. x k = 0)\}$ 
  proof
  show ?A  $\subseteq$   $\{\lambda n. 1\} \cup \{x. (\forall n. x n \in \{0, 1\}) \wedge (\forall n. \exists k \geq n. x k = 0)\}$ 
  proof
  fix  $x$ 
  assume  $x \in ?A$ 
  then obtain  $r$  where  $hr: r \in \{0..1\}$   $x = \text{to-Cantor-from-01 } r$ 
  by auto
  then consider  $r = 1 \mid r \in \{0..<1\}$  by fastforce
  thus  $x \in \{\lambda n. 1\} \cup \{x. (\forall n. x n \in \{0,1\}) \wedge (\forall n. \exists k \geq n. x k = 0)\}$ 
  proof cases
  case 1
  then show ?thesis
  by(simp add: hr(2) to-Cantor-from-01-def)
  next
  case 2
  from to-Cantor-from-01-exist0[OF this] to-Cantor-from-01-image'
  show ?thesis by(auto simp: hr(2))
  qed
qed

```

```

next
show  $\{\lambda n. 1\} \cup \{x. (\forall n. x n \in \{0, 1\}) \wedge (\forall n. \exists k \geq n. x k = 0)\} \subseteq ?A$ 
proof
  fix  $x :: nat \Rightarrow real$ 
  assume  $x \in \{\lambda n. 1\} \cup \{x. (\forall n. x n \in \{0, 1\}) \wedge (\forall n. \exists k \geq n. x k = 0)\}$ 
  then consider  $x = (\lambda n. 1) \mid (\forall n. x n \in \{0, 1\}) \wedge (\forall n. \exists k \geq n. x k =$ 
0)
    by auto
    thus  $x \in ?A$ 
  proof cases
    case 1
    then show ?thesis
    by(auto intro!: image-eqI[where x=1] simp: to-Cantor-from-01-def)
  next
    case 2
    hence  $\bigwedge n. 0 \leq x n \wedge n. x n \leq 1$ 
  by (metis dual-order.refl empty-iff insert-iff zero-less-one-class.zero-le-one)+
  with 2 to-Cantor-from-01-if-exist0[of x] nsum-of-r-leq[of 1/2 x 1 1 0]
  show ?thesis
  by(auto intro!: image-eqI[where x= $\sum n. (1 / 2) ^ Suc n * x n$ ])
  qed
  qed
  qed
  have  $(\Pi_E i \in UNIV. \{0, 1\}) - ?A = \{x. (\forall n. x n \in \{0, 1\}) \wedge (\exists n. \forall k \geq n. x k = 1)\} - \{\lambda n. 1\}$ 
proof
  show  $(\Pi_E i \in UNIV. \{0, 1\}) - ?A \subseteq \{x. (\forall n. x n \in \{0, 1\}) \wedge (\exists n. \forall k \geq n. x k = 1)\} - \{\lambda n. 1\}$ 
proof
  fix  $x :: nat \Rightarrow real$ 
  assume  $x \in (\Pi_E i \in UNIV. \{0, 1\}) - ?A$ 
  then have  $\forall n. x n \in \{0, 1\} \wedge (\forall n. \exists k \geq n. x k = 0) \wedge x \neq (\lambda n. 1)$ 
using eq1 by blast+
  thus  $x \in \{x. (\forall n. x n \in \{0, 1\}) \wedge (\exists n. \forall k \geq n. x k = 1)\} - \{\lambda n. 1\}$ 
by blast
  qed
next
show  $(\Pi_E i \in UNIV. \{0, 1\}) - ?A \supseteq \{x. (\forall n. x n \in \{0, 1\}) \wedge (\exists n. \forall k \geq n. x k = 1)\} - \{\lambda n. 1\}$ 
proof
  fix  $x :: nat \Rightarrow real$ 
  assume  $h: x \in \{x. (\forall n. x n \in \{0, 1\}) \wedge (\exists n. \forall k \geq n. x k = 1)\} - \{\lambda n. 1\}$ 
  then have  $\forall n. x n \in \{0, 1\} \wedge \exists n. \forall k \geq n. x k = 1 \wedge x \neq (\lambda n. 1)$ 
by blast+
  hence  $\neg (\forall n. \exists k \geq n. x k = 0)$ 
by fastforce
  with  $\langle \forall n. x n \in \{0, 1\} \rangle \langle x \neq (\lambda n. 1) \rangle$ 
show  $x \in (\Pi_E i \in UNIV. \{0, 1\}) - ?A$ 

```

```

    using eq1 by blast
  qed
  qed
  also have ... = (UNIV)) - {λn. 1}
    by blast
  also have ... ∈ sets Cantor-space (is ?B ∈ -)
  proof -
    have countable ?B
    proof -
      have countable {x :: nat ⇒ real. (∀ n. x n = 0 ∨ x n = 1) ∧ (∀ k ≥ m.
x k = 1)} for m :: nat
      proof -
        let ?C = {x::nat ⇒ real. (∀ n. x n = 0 ∨ x n = 1) ∧ (∀ k ≥ m. x k =
1)}
        define g where g = (λ(x::nat ⇒ real) n. if n < m then x n else
undefined)
        have 1: g ' ?C = (ΠE i ∈ {..m}. {0,1})
        proof(standard; standard)
          fix x
          assume x ∈ g ' ?C
          then show x ∈ (ΠE i ∈ {..m}. {0,1})
            by(auto simp: g-def PiE-def extensional-def)
        next
          fix x
          assume h: x ∈ (ΠE i ∈ {..m}. {0,1::real})
          then have x = g (λn. if n < m then x n else 1)
            by(auto simp add: g-def PiE-def extensional-def)
          moreover have (λn. if n < m then x n else 1) ∈ ?C
            using h by auto
          ultimately show x ∈ g ' ?C
            by auto
        qed
        have 2: inj-on g ?C
      proof
        fix x y
        assume hxyg: x ∈ ?C y : ?C g x = g y
        show x = y
        proof
          fix n
          consider n < m | m ≤ n by fastforce
          thus x n = y n
        proof cases
          case 1
          then show ?thesis
            using fun-cong[OF hxyg(3), of n] by(simp add: g-def)
        next
          case 2
          then show ?thesis

```

```

        using hxyg(1,2) by auto
      qed
    qed
  qed
  show countable {x::nat ⇒ real. (∀ n. x n = 0 ∨ x n = 1) ∧ (∀ k ≥ m.
x k = 1)}
    by(rule countable-image-inj-on[OF - 2]) (auto intro!: countable-PiE
simp: 1)
  qed
  thus ?thesis
    by auto
  qed
  moreover have ?B ⊆ space Cantor-space
    by(auto simp: space-Cantor-space)
  ultimately show ?thesis
    using Cantor-space-standard-ne by(simp add: standard-borel.countable-sets
standard-borel-ne-def)
  qed
  finally show ?thesis .
  qed
  moreover have space Cantor-space - ((ΠE i ∈ UNIV. {0,1}) - ?A) = ?A
    using ⟨?A ⊆ space Cantor-space⟩ space-Cantor-space by blast
  ultimately show ?thesis
    using sets.compl-sets[OF comple-sets] by auto
  qed
  have to-Cantor-from-01 ‹A = f -‹ A ∩ to-Cantor-from-01 ‹(space (restrict-space
borel {0..1}))››
  proof
    show to-Cantor-from-01 ‹A ⊆ f -‹ A ∩ to-Cantor-from-01 ‹space
(restrict-space borel {0..1})››
    proof
      fix x
      assume x ∈ to-Cantor-from-01 ‹A
      then obtain a where ha:a ∈ A x = to-Cantor-from-01 a by auto
      hence a ∈ {0..1}
        using sets.sets-into-space[OF assms] by auto
      have f x = a
        using to-Cantor-from-sum[OF ‹a ∈ {0..1}›] by(simp add: f-def ha(2))
      thus x ∈ f -‹ A ∩ to-Cantor-from-01 ‹space (restrict-space borel {0..1})››
        using sets.sets-into-space[OF assms] ha by auto
    qed
  next
    show to-Cantor-from-01 ‹A ⊇ f -‹ A ∩ to-Cantor-from-01 ‹space
(restrict-space borel {0..1})››
    proof
      fix x
      assume h:x ∈ f -‹ A ∩ to-Cantor-from-01 ‹space (restrict-space borel
{0..1})››
      then obtain r where r ∈ {0..1} x = to-Cantor-from-01 r

```

```

    by auto
    from h have f x ∈ A
    by simp
    hence to-Cantor-from-01 (f x) = x
    using to-Cantor-from-01-sum-of-to-Cantor-from-01[OF ‹r ∈ {0..1}›]
    by(simp add: f-def ‹x = to-Cantor-from-01 r›)
    with ‹f x ∈ A›
    show x ∈ to-Cantor-from-01 ‹ A
    by (simp add: rev-image-eqI)
  qed
  qed
  also have ... ∈ sets Cantor-space
  proof -
    have f - ‹ A ∩ space Cantor-space ∩ to-Cantor-from-01 ‹ space (restrict-space
    borel {0..1}) = f - ‹ A ∩ to-Cantor-from-01 ‹ (space (restrict-space borel {0..1}))
    using to-Cantor-from-01-image sets.sets-into-space[OF assms,simplified]
  by auto
    thus ?thesis
    using sets.Int[OF measurable-sets[OF f-meas assms] image-sets]
    by fastforce
  qed
  finally show ?thesis .
  qed
  show ?thesis
    using Schroeder-Bernstein-measurable[OF Cantor-to-01-measurable Can-
    tor-to-01-preserves-sets Cantor-to-01-inj to-Cantor-from-01-measurable to-Cantor-from-01-preserves-sets
    to-Cantor-from-01-inj]
    by(simp add: measurable-isomorphic-def)
  qed
  have 1:Cantor-space measurable-isomorphic (ΠM (i::nat,j::nat)∈ UNIV × UNIV.
  restrict-space borel {0,1::real})
  unfolding Cantor-space-def
  by(auto intro!: measurable-isomorphic-sym[OF countable-infinite-isomorphisc-to-nat-index]
  simp: split-beta' finite-prod)
  have 2:(ΠM (i::nat,j::nat)∈ UNIV × UNIV. restrict-space borel {0,1::real})
  measurable-isomorphic (ΠM (i::nat)∈ UNIV. Cantor-space)
  unfolding Cantor-space-def by(rule measurable-isomorphic-sym[OF PiM-PiM-isomorphic-to-PiM])
  have 3:(ΠM (i::nat)∈ UNIV. Cantor-space) measurable-isomorphic Hilbert-cube
  unfolding Hilbert-cube-def by(rule measurable-isomorphic-lift-product[OF Can-
  tor-space-isomorphic-to-01closed])
  show ?thesis
  by(rule measurable-isomorphic-trans[OF measurable-isomorphic-trans[OF 1 2]
  3])
  qed

```

3.3 Final Results

lemma(in standard-borel) embedding-into-Hilbert-cube:

$\exists A \in \text{sets Hilbert-cube. } M \text{ measurable-isomorphic (restrict-space Hilbert-cube } A)$

proof –
obtain S **where** S :Polish-space S sets (borel-of S) = sets M
using Polish-space **by** blast
obtain A **where** A :gdelta-in Hilbert-cube-topology A S homeomorphic-space subtopology Hilbert-cube-topology A
using embedding-into-Hilbert-cube-gdelta-in[OF $S(1)$] **by** blast
show ?thesis
using borel-of-gdelta-in[OF $A(1)$] homeomorphic-space-measurable-isomorphic[OF $A(2)$] measurable-isomorphic-sets-cong[OF $S(2)$,of borel-of (subtopology Hilbert-cube-topology A) restrict-space Hilbert-cube A] Hilbert-cube-borel sets-restrict-space-cong[OF Hilbert-cube-borel]
by(auto intro!: bexI[**where** $x=A$] simp: borel-of-subtopology)
qed

lemma(in standard-borel) embedding-from-Cantor-space:

assumes uncountable (space M)
shows $\exists A \in$ sets M . Cantor-space measurable-isomorphic (restrict-space M A)

proof –
obtain S **where** S :Polish-space S sets (borel-of S) = sets M
using Polish-space **by** blast
then obtain A **where** A :gdelta-in S A Cantor-space-topology homeomorphic-space subtopology S A
using embedding-from-Cantor-space[of S] assms sets-eq-imp-space-eq[OF $S(2)$] **by**(auto simp: space-borel-of)
show ?thesis
using borel-of-gdelta-in[OF $A(1)$] $S(2)$ homeomorphic-space-measurable-isomorphic[OF $A(2)$] measurable-isomorphic-sets-cong[OF Cantor-space-borel restrict-space-sets-cong[OF refl $S(2)$],of A]
by(auto intro!: bexI[**where** $x=A$] simp: borel-of-subtopology)
qed

corollary(in standard-borel) uncountable-isomorphic-to-Hilbert-cube:

assumes uncountable (space M)
shows Hilbert-cube measurable-isomorphic M

proof –
obtain A B **where** AB :
 M measurable-isomorphic (restrict-space Hilbert-cube A) Cantor-space measurable-isomorphic (restrict-space M B)
 $A \in$ sets Hilbert-cube $B \in$ sets M
using embedding-into-Hilbert-cube embedding-from-Cantor-space[OF assms] **by** auto
show ?thesis
by(rule measurable-isomorphic-antisym[OF AB measurable-isomorphic-sym[OF Cantor-space-isomorphic-to-Hilbert-cube]])
qed

corollary(in standard-borel) uncountable-isomorphic-to-real:

assumes uncountable (space M)
shows M measurable-isomorphic (borel :: real measure)

proof –


```

interpret r: standard-borel-ne borel :: real measure
  by simp
show ?thesis
  by(auto intro!: measurable-isomorphic-trans[OF measurable-isomorphic-sym[OF
uncountable-isomorphic-to-Hilbert-cube[OF assms]] r.uncountable-isomorphic-to-Hilbert-cube
simp: uncountable-UNIV-real)
qed

lemma(in standard-borel isomorphic-subset-real:
  assumes A ∈ sets (borel :: real measure) uncountable A
  obtains B where B ∈ sets borel B ⊆ A M measurable-isomorphic restrict-space
borel B
proof(cases countable (space M))
  assume count:countable (space M)
  have ∃ B⊆A. space M ≈ B
  proof(cases finite (space M))
    assume fin:finite (space M)
    then obtain B where B:card B = card (space M) finite B B ⊆ A
    by (meson assms(2) countable-finite infinite-arbitrarily-large)
    thus ?thesis
    using fin by(auto intro!: exI[where x=B] simp: eqpoll-iff-card)
  next
    assume inf:infinite (space M)
    obtain B where B: B ⊆ A countable B infinite B
    using assms(2) countable-finite infinite-countable-subset' that by auto
    thus ?thesis
    using bij-betw-from-nat-into[OF count inf] bij-betw-from-nat-into[OF B(2,3)]
    by (meson eqpoll-def eqpoll-sym eqpoll-trans)
  qed
then obtain B where B:B ⊆ A space M ≈ B countable B
  by (metis countable-reqpoll eqpoll-sym count)
then obtain f where f:bij-betw f (space M) B
  using eqpoll-def by blast
have 1:C ∈ sets borel if C:C ⊆ B for C
proof –
  have C = (⋃ c∈C. {c})
  by auto
  also have ... ∈ sets borel
  using B C by(intro sets.countable-UN') (auto simp: countable-subset)
  finally show ?thesis .
qed
have 2:sets M = sets (count-space (space M))
  by (simp add: countable-discrete-space count)
have 3:sets (restrict-space borel B) = sets (count-space B)
  using 1 by(auto simp: sets-restrict-space)
have [simp]:measurable M (restrict-space borel B) = measurable (count-space
(space M)) (count-space B)
  measurable (restrict-space borel B) M = measurable (count-space B) (count-space
(space M))

```

using 2 3 **by**(*auto intro!*: *measurable-cong-sets*)
have *M measurable-isomorphic restrict-space borel B*
using *bij-betw-the-inv-into[OF f] f* **by**(*auto simp: measurable-isomorphic-def measurable-isomorphic-map-def space-restrict-space intro!: exI[where x=f] dest: bij-betwE*)
with 1 *B that show ?thesis*
by *blast*
next
assume *uncountable (space M)*
then have *M measurable-isomorphic (borel :: real measure)*
using *uncountable-isomorphic-to-real* **by** *blast*
moreover have *restrict-space borel A measurable-isomorphic (borel :: real measure)*
by(*auto intro!: standard-borel.uncountable-isomorphic-to-real standard-borel.standard-borel-restrict-space[OF standard-borel-ne.standard-borel] simp: assms space-restrict-space*)
ultimately have *M measurable-isomorphic restrict-space borel A*
using *measurable-isomorphic-sym measurable-isomorphic-trans* **by** *blast*
with *assms(1) that show ?thesis*
by *blast*
qed

lemma(*in standard-borel*) *countable-isomorphic-to-subset-real*:
assumes *countable (space M)*
obtains *A :: real set*
where *countable A A ∈ sets borel M measurable-isomorphic restrict-space borel A*
proof –
obtain *A :: real set where A:A ∈ sets borel M measurable-isomorphic restrict-space borel A*
using *isomorphic-subset-real[of UNIV] uncountable-UNIV-real* **by** *auto*
moreover have *countable A*
using *measurable-isomorphic-cardinality-eq[OF A(2)] assms(1)*
by(*simp add: space-restrict-space countable-eqpoll[OF - eqpoll-sym]*)
ultimately show *?thesis*
using *that* **by** *blast*
qed

theorem *Borel-isomorphism-theorem*:
assumes *standard-borel M standard-borel N*
shows *space M ≈ space N ⟷ M measurable-isomorphic N*
proof
assume *h:space M ≈ space N*
interpret *M: standard-borel M* **by** *fact*
interpret *N: standard-borel N* **by** *fact*
consider *countable (space M) countable (space N) | uncountable (space M) uncountable (space N)*
by (*meson countable-eqpoll eqpoll-sym h*)
thus *M measurable-isomorphic N*
proof *cases*

```

case 1
then have  $2$ :sets  $M = \text{sets } (\text{count-space } (\text{space } M)) \text{ sets } N = \text{sets } (\text{count-space } (\text{space } N))$ 
by (simp-all add:  $M$ .countable-discrete-space  $N$ .countable-discrete-space)
show ?thesis
by(simp add: measurable-isomorphic-sets-cong[OF  $2$ ] measurable-isomorphic-count-spaces
h)
next
case 2
then have  $M$  measurable-isomorphic (borel :: real measure)  $N$  measurable-isomorphic
(borel :: real measure)
by(simp-all add:  $M$ .uncountable-isomorphic-to-real  $N$ .uncountable-isomorphic-to-real)
thus ?thesis
using measurable-isomorphic-sym measurable-isomorphic-trans by blast
qed
qed(rule measurable-isomorphic-cardinality-eq)

```

definition *to-real-on* :: 'a measure \Rightarrow 'a \Rightarrow real **where**
to-real-on $M \equiv (\text{if uncountable } (\text{space } M) \text{ then } (\text{SOME } f. \text{measurable-isomorphic-map } M \text{ (borel :: real measure) } f) \text{ else } (\text{real } \circ \text{to-nat-on } (\text{space } M)))$

definition *from-real-into* :: 'a measure \Rightarrow real \Rightarrow 'a **where**
from-real-into $M \equiv (\text{if uncountable } (\text{space } M) \text{ then the-inv-into } (\text{space } M) \text{ (to-real-on } M) \text{ else } (\lambda r. \text{from-nat-into } (\text{space } M) \text{ (nat } \lfloor r \rfloor)))$

context *standard-borel*
begin

abbreviation *to-real* \equiv *to-real-on* M
abbreviation *from-real* \equiv *from-real-into* M

lemma *to-real-def-countable*:
assumes *countable* (space M)
shows *to-real* = ($\lambda r. \text{real } (\text{to-nat-on } (\text{space } M) r)$)
using *assms* **by**(auto simp: *to-real-on-def*)

lemma *from-real-def-countable*:
assumes *countable* (space M)
shows *from-real* = ($\lambda r. \text{from-nat-into } (\text{space } M) \text{ (nat } \lfloor r \rfloor)$)
using *assms* **by**(simp add: *from-real-into-def*)

lemma *from-real-to-real[simp]*:
assumes $x \in \text{space } M$
shows *from-real* (to-real x) = x
proof –
have [simp]: $\text{space } M \neq \{\}$
using *assms* **by** auto
consider *countable* (space M) | *uncountable* (space M) **by** auto
then show ?thesis

```

proof cases
  case 1
  then show ?thesis
    by(simp add: to-real-def-countable from-real-def-countable assms)
  next
  case 2
  then obtain f where f: measurable-isomorphic-map M (borel :: real measure)
f
  using uncountable-isomorphic-to-real by(auto simp: measurable-isomorphic-def)
  have 1:to-real = Eps (measurable-isomorphic-map M borel) from-real = the-inv-into
(space M) (Eps (measurable-isomorphic-map M borel))
  by(simp-all add: to-real-on-def 2 from-real-into-def)
  show ?thesis
  unfolding 1
  by(rule someI2[of measurable-isomorphic-map M (borel :: real measure) f, OF
f])
  (meson assms bij-betw-imp-inj-on measurable-isomorphic-map-def the-inv-into-f-f)
  qed
qed

```

```

lemma to-real-measurable[measurable]:
  to-real ∈ M →M borel
proof(cases countable (space M))
  case 1:True
  then have sets M = Pow (space M)
  by(rule countable-discrete-space)
  then show ?thesis
  by(simp add: to-real-def-countable 1 borel-measurableI-le)
next
  case 1:False
  then obtain f where f: measurable-isomorphic-map M (borel :: real measure) f
  using uncountable-isomorphic-to-real by(auto simp: measurable-isomorphic-def)
  have 2:to-real = Eps (measurable-isomorphic-map M borel)
  by(simp add: to-real-on-def 1 from-real-into-def)
  show ?thesis
  unfolding 2
  by(rule someI2[of measurable-isomorphic-map M (borel :: real measure) f, OF
f],simp add: measurable-isomorphic-map-def)
qed

```

```

lemma from-real-measurable':
  assumes space M ≠ {}
  shows from-real ∈ borel →M M
proof(cases countable (space M))
  case 1:True
  then have 2:sets M = Pow (space M)
  by(rule countable-discrete-space)
  have [measurable]:from-nat-into (space M) ∈ count-space UNIV →M M
  using from-nat-into[OF assms] by auto

```

```

show ?thesis
  by(simp add: from-real-def-countable 1 borel-measurableI-le)
next
  case 2:False
  then obtain f where f: measurable-isomorphic-map M (borel :: real measure) f
  using uncountable-isomorphic-to-real by(auto simp: measurable-isomorphic-def)
  have 1: from-real = the-inv-into (space M) (Eps (measurable-isomorphic-map M
borel))
  by(simp add: to-real-on-def 2 from-real-into-def)
  show ?thesis
  unfolding 1
  by(rule someI2[of measurable-isomorphic-map M (borel :: real measure) f, OF
f], simp add: measurable-isomorphic-map-def)
qed

lemma to-real-from-real:
  assumes uncountable (space M)
  shows to-real (from-real r) = r
proof –
  obtain f where f: measurable-isomorphic-map M (borel :: real measure) f
  using assms uncountable-isomorphic-to-real by(auto simp: measurable-isomorphic-def)
  have 1: to-real = Eps (measurable-isomorphic-map M borel) from-real = the-inv-into
(space M) (Eps (measurable-isomorphic-map M borel))
  by(simp-all add: to-real-on-def assms from-real-into-def)
  show ?thesis
  unfolding 1
  by(rule someI2[of measurable-isomorphic-map M (borel :: real measure) f, OF
f])
  (metis UNIV-I f-the-inv-into-f-bij-betw measurable-isomorphic-map-def space-borel)
qed

end

lemma(in standard-borel-ne) from-real-measurable[measurable]: from-real ∈ borel
→M M
  by(simp add: from-real-measurable' space-ne)

end

```

References

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