

A Complete Proof of the Robbins Conjecture

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Abstract

The document gives a formalization of the proof of the Robbins conjecture, following A. Mann, *A Complete Proof of the Robbins Conjecture*, 2003.

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1 Robbins Conjecture

```
theory Robbins-Conjecture
imports Main
begin
```

The document gives a formalization of the proof of the Robbins conjecture, following A. Mann, *A Complete Proof of the Robbins Conjecture*, 2003, DOI 10.1.1.6.7838

2 Axiom Systems

The following presents several axiom systems that shall be under study.

The first axiom sets common systems that underly all of the systems we shall be looking at.

The second system is a reformulation of Boolean algebra. We shall follow pages 7–8 in S. Koppelberg. *General Theory of Boolean Algebras*, Volume 1 of *Handbook of Boolean Algebras*. North Holland, 1989. Note that our formulation deviates slightly from this, as we only provide one distribution axiom, as the dual is redundant.

The third system is Huntington’s algebra and the fourth system is Robbins’ algebra.

Apart from the common system, all of these systems are demonstrated to be equivalent to the library formulation of Boolean algebra, under appropriate interpretation.

2.1 Common Algebras

```

class common-algebra = uminus +
  fixes inf :: 'a ⇒ 'a ⇒ 'a (infixl □ 70)
  fixes sup :: 'a ⇒ 'a ⇒ 'a (infixl ⊔ 65)
  fixes bot :: 'a (⊥)
  fixes top :: 'a (⊤)
  assumes sup-assoc: x ⊔ (y ⊔ z) = (x ⊔ y) ⊔ z
  assumes sup-comm: x ⊔ y = y ⊔ x

context common-algebra begin

definition less-eq :: 'a ⇒ 'a ⇒ bool (infix ⊑ 50) where
  x ⊑ y = (x ⊔ y = y)
definition less :: 'a ⇒ 'a ⇒ bool (infix ⊒ 50) where
  x ⊒ y = (x ⊑ y ∧ ¬ y ⊑ x)
definition minus :: 'a ⇒ 'a ⇒ 'a (infixl − 65) where
  minus x y = (x ⊓ − y)

definition secret-object1 :: 'a (ι) where
  ι = (SOME x. True)

end

class ext-common-algebra = common-algebra +
  assumes inf-eq: x ⊓ y = −(− x ⊔ − y)
  assumes top-eq: ⊤ = ι ⊔ − ι
  assumes bot-eq: ⊥ = −(ι ⊔ − ι)

```

2.2 Boolean Algebra

```

class boolean-algebra-II =
  common-algebra +
  assumes inf-comm: x ⊓ y = y ⊓ x

```

```

assumes inf-assoc:  $x \sqcap (y \sqcap z) = (x \sqcap y) \sqcap z$ 
assumes sup-absorb:  $x \sqcup (x \sqcap y) = x$ 
assumes inf-absorb:  $x \sqcap (x \sqcup y) = x$ 
assumes sup-inf-distrib1:  $x \sqcup y \sqcap z = (x \sqcup y) \sqcap (x \sqcup z)$ 
assumes sup-compl:  $x \sqcup -x = \top$ 
assumes inf-compl:  $x \sqcap -x = \perp$ 

```

2.3 Huntington's Algebra

```

class huntington-algebra = ext-common-algebra +
  assumes huntington:  $-(-x \sqcup -y) \sqcup -(x \sqcup y) = x$ 

```

2.4 Robbins' Algebra

```

class robbins-algebra = ext-common-algebra +
  assumes robbins:  $-(-(x \sqcup y) \sqcup -(x \sqcup -y)) = x$ 

```

3 Equivalence

With our axiom systems defined, we turn to providing equivalence results between them.

We shall begin by illustrating equivalence for our formulation and the library formulation of Boolean algebra.

3.1 Boolean Algebra

The following provides the canonical definitions for order and relative complementation for Boolean algebras. These are necessary since the Boolean algebras presented in the Isabelle/HOL library have a lot of structure, while our formulation is considerably simpler.

Since our formulation of Boolean algebras is considerably simple, it is easy to show that the library instantiates our axioms.

```

context boolean-algebra-II begin

```

```

lemma boolean-II-is-boolean:

```

```

  class.boolean-algebra minus uminus ( $\sqcap$ ) ( $\sqcup$ ) ( $\sqcap$ ) ( $\sqcup$ )  $\perp$   $\top$ 
  <proof>

```

```

end

```

```

context boolean-algebra begin

```

```

lemma boolean-is-boolean-II:

```

```

  class.boolean-algebra-II uminus inf sup bot top
  <proof>

```

```

end

```

3.2 Huntington Algebra

We shall illustrate here that all Boolean algebra using our formulation are Huntington algebras, and illustrate that every Huntington algebra may be interpreted as a Boolean algebra.

Since the Isabelle/HOL library has good automation, it is convenient to first show that the library instances Huntington algebras to exploit previous results, and then use our previously derived correspondence.

context *boolean-algebra* **begin**

lemma *boolean-is-huntington*:

class.huntington-algebra uminus inf sup bot top
<proof>

end

context *boolean-algebra-II* **begin**

lemma *boolean-II-is-huntington*:

class.huntington-algebra uminus (\sqcap) (\sqcup) \perp \top
<proof>

end

context *huntington-algebra* **begin**

lemma *huntington-id*: $x \sqcup -x = -x \sqcup -(-x)$

<proof>

lemma *dbl-neg*: $-(-x) = x$

<proof>

lemma *towards-sup-compl*: $x \sqcup -x = y \sqcup -y$

<proof>

lemma *sup-compl*: $x \sqcup -x = \top$

<proof>

lemma *towards-inf-compl*: $x \sqcap -x = y \sqcap -y$

<proof>

lemma *inf-compl*: $x \sqcap -x = \perp$

<proof>

lemma *towards-idem*: $\perp = \perp \sqcup \perp$

<proof>

lemma *sup-ident*: $x \sqcup \perp = x$

<proof>

lemma *inf-ident*: $x \sqcap \top = x$
<proof>

lemma *sup-idem*: $x \sqcup x = x$
<proof>

lemma *inf-idem*: $x \sqcap x = x$
<proof>

lemma *sup-nil*: $x \sqcup \top = \top$
<proof>

lemma *inf-nil*: $x \sqcap \perp = \perp$
<proof>

lemma *sup-absorb*: $x \sqcup x \sqcap y = x$
<proof>

lemma *inf-absorb*: $x \sqcap (x \sqcup y) = x$
<proof>

lemma *partition*: $x \sqcap y \sqcup x \sqcap \neg y = x$
<proof>

lemma *demorgans1*: $\neg(x \sqcap y) = \neg x \sqcup \neg y$
<proof>

lemma *demorgans2*: $\neg(x \sqcup y) = \neg x \sqcap \neg y$
<proof>

lemma *inf-comm*: $x \sqcap y = y \sqcap x$
<proof>

lemma *inf-assoc*: $x \sqcap (y \sqcap z) = x \sqcap y \sqcap z$
<proof>

lemma *inf-sup-distrib1*: $x \sqcap (y \sqcup z) = (x \sqcap y) \sqcup (x \sqcap z)$
<proof>

lemma *sup-inf-distrib1*:
 $x \sqcup (y \sqcap z) = (x \sqcup y) \sqcap (x \sqcup z)$
<proof>

lemma *huntington-is-boolean-II*:
class.boolean-algebra-II *uminus* (\sqcap) (\sqcup) \perp \top
<proof>

lemma *huntington-is-boolean*:
class.boolean-algebra *minus* *uminus* (\sqcap) (\sqcup) (\sqsubseteq) (\sqsupseteq) (\sqcup) \perp \top

<proof>
end

3.3 Robbins' Algebra

context *boolean-algebra* **begin**
lemma *boolean-is-robbins*:
 class.robbins-algebra uminus inf sup bot top
<proof>
end

context *boolean-algebra-II* **begin**
lemma *boolean-II-is-robbins*:
 class.robbins-algebra uminus inf sup bot top
<proof>
end

context *huntington-algebra* **begin**
lemma *huntington-is-robbins*:
 class.robbins-algebra uminus inf sup bot top
<proof>
end

Before diving into the proof that the Robbins algebra is Boolean, we shall present some shorthand machinery

context *common-algebra* **begin**

primrec *copy* :: *nat* \Rightarrow *'a* \Rightarrow *'a* (**infix** \otimes 80)
where
 copy-0: $0 \otimes x = x$
 | *copy-Suc*: $(\text{Suc } k) \otimes x = (k \otimes x) \sqcup x$

no-notation
 Product-Type.Times (**infixr** \times 80)

primrec *copy* :: *nat* \Rightarrow *'a* \Rightarrow *'a* (**infix** \times 85)
where
 $0 \times x = x$
 | $(\text{Suc } k) \times x = k \otimes x$

lemma *one*: $1 \times x = x$
<proof>

lemma *two*: $2 \times x = x \sqcup x$
<proof>

lemma three: $3 \times x = x \sqcup x \sqcup x$
<proof>

lemma four: $4 \times x = x \sqcup x \sqcup x \sqcup x$
<proof>

lemma five: $5 \times x = x \sqcup x \sqcup x \sqcup x \sqcup x$
<proof>

lemma six: $6 \times x = x \sqcup x \sqcup x \sqcup x \sqcup x \sqcup x$
<proof>

lemma cotyp-distrib: $k \otimes (x \sqcup y) = (k \otimes x) \sqcup (k \otimes y)$
<proof>

corollary copy-distrib: $k \times (x \sqcup y) = (k \times x) \sqcup (k \times y)$
<proof>

lemma cotyp-arith: $(k + l + 1) \otimes x = (k \otimes x) \sqcup (l \otimes x)$
<proof>

lemma copy-arith:

assumes $k \neq 0$ **and** $l \neq 0$

shows $(k + l) \times x = (k \times x) \sqcup (l \times x)$

<proof>

end

The theorem asserting all Robbins algebras are Boolean comes in 6 movements.

First: The Winker identity is proved.

Second: Idempotence for a particular object is proved. Note that falsum is defined in terms of this object.

Third: An identity law for falsum is derived.

Fourth: Idempotence for supremum is derived.

Fifth: The double negation law is proven

Sixth: Robbin's algebras are proven to be Huntington Algebras.

context robbins-algebra begin

definition secret-object2 :: 'a (α) **where**

$\alpha = -(-(\iota \sqcup \iota \sqcup \iota) \sqcup \iota)$

definition secret-object3 :: 'a (β) **where**

$\beta = \iota \sqcup \iota$

definition secret-object4 :: 'a (δ) **where**

$\delta = \beta \sqcup (-(\alpha \sqcup -\beta) \sqcup -(\alpha \sqcup -\beta))$

definition *secret-object5* :: 'a (γ) **where**

$$\gamma = \delta \sqcup -(\delta \sqcup -\delta)$$

definition *winker-object* :: 'a (ϱ) **where**

$$\varrho = \gamma \sqcup \gamma \sqcup \gamma$$

definition *fake-bot* :: 'a ($\perp\perp$) **where**

$$\perp\perp = -(\varrho \sqcup -\varrho)$$

lemma *robbins2*: $y = -(-(-x \sqcup y) \sqcup -(x \sqcup y))$

<proof>

lemma *mann0*: $-(x \sqcup y) = -(-(-(-x \sqcup y) \sqcup -x \sqcup y) \sqcup y)$

<proof>

lemma *mann1*: $-(-x \sqcup y) = -(-(-(-x \sqcup y) \sqcup x \sqcup y) \sqcup y)$

<proof>

lemma *mann2*: $y = -(-(-(-x \sqcup y) \sqcup x \sqcup y \sqcup y) \sqcup -(x \sqcup y))$

<proof>

lemma *mann3*: $z = -(-(-(-(-x \sqcup y) \sqcup x \sqcup y \sqcup y) \sqcup -(x \sqcup y) \sqcup z) \sqcup -(y \sqcup z))$

<proof>

lemma *mann4*: $-(y \sqcup z) =$

$$-(-(-(-(-x \sqcup y) \sqcup x \sqcup y \sqcup y) \sqcup -(x \sqcup y) \sqcup -(y \sqcup z) \sqcup z) \sqcup z)$$

<proof>

lemma *mann5*: $u =$

$$\begin{aligned} & -(-(-(-(-(-x \sqcup y) \sqcup x \sqcup y \sqcup y) \sqcup \\ & \quad -(-x \sqcup y) \sqcup -(y \sqcup z) \sqcup z) \sqcup z \sqcup u) \sqcup \\ & \quad -(-(y \sqcup z) \sqcup u)) \end{aligned}$$

<proof>

lemma *mann6*:

$$-(-\ 3 \times x \sqcup x) = -(-(-(-\ 3 \times x \sqcup x) \sqcup -\ 3 \times x) \sqcup -(-(-\ 3 \times x \sqcup x) \sqcup 5 \times x))$$

<proof>

lemma *mann7*:

$$-\ 3 \times x = -(-(-\ 3 \times x \sqcup x) \sqcup 5 \times x)$$

<proof>

lemma *mann8*:

$$-(-\ 3 \times x \sqcup x) \sqcup 2 \times x = -(-(-(-\ 3 \times x \sqcup x) \sqcup -\ 3 \times x \sqcup 2 \times x) \sqcup -\ 3 \times x)$$

(**is** ?lhs = ?rhs)

<proof>

lemma mann9: $x = -(-(-\mathcal{I} \times x \sqcup x) \sqcup -\mathcal{I} \times x)$
<proof>

lemma mann10: $y = -(-(-(-\mathcal{I} \times x \sqcup x) \sqcup -\mathcal{I} \times x \sqcup y) \sqcup -(x \sqcup y))$
<proof>

theorem mann: $2 \times x = -(-\mathcal{I} \times x \sqcup x) \sqcup 2 \times x$
<proof>

corollary winkerr: $\alpha \sqcup \beta = \beta$
<proof>

corollary winker: $\beta \sqcup \alpha = \beta$
<proof>

corollary multi-winkerp: $\beta \sqcup k \otimes \alpha = \beta$
<proof>

corollary multi-winker: $\beta \sqcup k \times \alpha = \beta$
<proof>

lemma less-eq-introp:
 $-(x \sqcup -(y \sqcup z)) = -(x \sqcup y \sqcup -z) \implies y \sqsubseteq x$
<proof>

corollary less-eq-intro:
 $-(x \sqcup -(y \sqcup z)) = -(x \sqcup y \sqcup -z) \implies x \sqcup y = x$
<proof>

lemma eq-intro:
 $-(x \sqcup -(y \sqcup z)) = -(y \sqcup -(x \sqcup z)) \implies x = y$
<proof>

lemma copyp0:
 assumes $-(x \sqcup -y) = z$
 shows $-(x \sqcup -(y \sqcup k \otimes (x \sqcup z))) = z$
<proof>

lemma copyp1:
 assumes $-(-(x \sqcup -y) \sqcup -y) = x$
 shows $-(y \sqcup k \otimes (x \sqcup -(x \sqcup -y))) = -y$
<proof>

corollary copyp2:
 assumes $-(x \sqcup y) = -y$
 shows $-(y \sqcup k \otimes (x \sqcup -(x \sqcup -y))) = -y$

$\langle proof \rangle$

lemma two-threep:

assumes $-(2 \times x \sqcup y) = -y$

and $-(3 \times x \sqcup y) = -y$

shows $2 \times x \sqcup y = 3 \times x \sqcup y$

$\langle proof \rangle$

lemma two-three:

assumes $-(x \sqcup y) = -y \vee -(-(x \sqcup -y) \sqcup -y) = x$

shows $y \sqcup 2 \times (x \sqcup -(x \sqcup -y)) = y \sqcup 3 \times (x \sqcup -(x \sqcup -y))$

(is $y \sqcup ?z2 = y \sqcup ?z3$)

$\langle proof \rangle$

lemma sup-idem: $\varrho \sqcup \varrho = \varrho$

$\langle proof \rangle$

lemma sup-ident: $x \sqcup \perp\perp = x$

$\langle proof \rangle$

lemma dbl-neg: $-(-x) = x$

$\langle proof \rangle$

theorem robbins-is-huntington:

class.huntington-algebra *uminus* (\sqcap) (\sqcup) \perp \top

$\langle proof \rangle$

theorem robbins-is-boolean-II:

class.boolean-algebra-II *uminus* (\sqcap) (\sqcup) \perp \top

$\langle proof \rangle$

theorem robbins-is-boolean:

class.boolean-algebra *minus* *uminus* (\sqcap) (\sqsubseteq) (\sqsupseteq) (\sqcup) \perp \top

$\langle proof \rangle$

end

no-notation *secret-object1* (ι)

and *secret-object2* (α)

and *secret-object3* (β)

and *secret-object4* (δ)

and *secret-object5* (γ)

and *winker-object* (ϱ)

```
and less-eq (infix  $\sqsubseteq$  50)
and less (infix  $\sqsubset$  50)
and inf (infixl  $\sqcap$  70)
and sup (infixl  $\sqcup$  65)
and top ( $\top$ )
and bot ( $\perp$ )
and copyp (infix  $\otimes$  80)
and copy (infix  $\times$  85)
```

notation

Product-Type.Times (**infixr** \times 80)

end