

Relational Disjoint-Set Forests

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Abstract

We give a simple relation-algebraic semantics of read and write operations on associative arrays. The array operations seamlessly integrate with assignments in the Hoare-logic library. Using relation algebras and Kleene algebras we verify the correctness of an array-based implementation of disjoint-set forests using the union-by-rank strategy and find operations with path compression, path splitting and path halving.

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1 Overview

Relation algebras and Kleene algebras have previously been used to reason about graphs and graph algorithms [2, 3, 4, 5, 9, 12, 15]. The operations of these algebras manipulate entire graphs, which is useful for specification but not directly intended for implementation. Low-level array access is a key ingredient for efficient algorithms [6]. We give a relation-algebraic semantics for such read/write access to associative arrays. This allows us to extend relation-algebraic verification methods to a lower level of more efficient implementations.

In this theory we focus on arrays with the same index and value sets, which can be modelled as homogeneous relations and therefore as elements of relation algebras and Kleene algebras [13, 17]. We implement and verify the correctness of disjoint-set forests with path compression strategies and union-by-rank [6, 8, 16].

In order to prepare this theory for future applications with weighted graphs, the verification uses Stone relation algebras, which have weaker axioms than relation algebras [10].

Section 2 contains the simple relation-algebraic semantics of associative array read and write and basic properties of these access operations. In Section 3 we give a Kleene-relation-algebraic semantics of disjoint-set forests. The make-set operation, find-set with path compression and the naive union-sets operation are implemented and verified in Section 4. Section 5 presents further results on disjoint-set forests and relational array access. The initialisation of disjoint-set forests, path halving and path splitting are implemented and verified in Section 6. In Section 7 we study relational Peano structures and implement and verify union-by-rank.

This Isabelle/HOL theory formally verifies results in [11] and in an extended version of that paper. Theorem numbers from the extended version of the paper are mentioned in the theories for reference. See the paper for further details and related work.

Several Isabelle/HOL theories are related to disjoint sets. The theory `HOL/Library/Disjoint_Sets.thy` contains results about partitions and sets of disjoint sets and does not consider their implementation. An implementation of disjoint-set forests with path compression and a size-based heuristic in the Imperative/HOL framework is verified in Archive of Formal Proofs entry [14]. Improved automation of this proof is considered in Archive of Formal Proofs entry [18]. These approaches are based on logical specifications whereas the present theory uses relation algebras and Kleene algebras.

theory *Disjoint-Set-Forests*

```

imports
  HOL-Hoare.Hoare-Logic
  Stone-Kleene-Relation-Algebras.Kleene-Relation-Algebras
begin

no-notation
  minus (infixl - 65) and
  trancl ((+) [1000] 999)

context p-algebra
begin

abbreviation minus :: 'a ⇒ 'a ⇒ 'a (infixl - 65)
  where  $x - y \equiv x \sqcap -y$ 

end

```

An arc in a Stone relation algebra corresponds to an atom in a relation algebra and represents a single edge in a graph. A point represents a set of nodes. A rectangle represents the Cartesian product of two sets of nodes [4].

```

context times-top
begin

abbreviation rectangle :: 'a ⇒ bool
  where  $rectangle\ x \equiv x * top * x = x$ 

end

context stone-relation-algebra
begin

lemma arc-rectangle:
   $arc\ x \implies rectangle\ x$ 
  using arc-top-arc by blast

```

2 Relation-Algebraic Semantics of Associative Array Access

The following two operations model updating array x at index y to value z , and reading the content of array x at index y , respectively. The read operation uses double brackets to avoid ambiguity with list syntax. The remainder of this section shows basic properties of these operations.

```

abbreviation rel-update :: 'a ⇒ 'a ⇒ 'a ⇒ 'a (([-⟶-]) [70, 65, 65] 61)
  where  $x[y \longrightarrow z] \equiv (y \sqcap z^T) \sqcup (-y \sqcap x)$ 

abbreviation rel-access :: 'a ⇒ 'a ⇒ 'a ((2-[-]) [70, 65] 65)

```

where $x[[y]] \equiv x^T * y$

Theorem 1.1

lemma *update-univalent*:

assumes *univalent* x

and *vector* y

and *injective* z

shows *univalent* $(x[y \mapsto z])$

proof –

have 1: *univalent* $(y \sqcap z^T)$

using *assms*(3) *inf-commute univalent-inf-closed* **by** *force*

have $(y \sqcap z^T)^T * (-y \sqcap x) = (y^T \sqcap z) * (-y \sqcap x)$

by (*simp add: conv-dist-inf*)

also have $\dots = z * (y \sqcap -y \sqcap x)$

by (*metis assms*(2) *covector-inf-comp-3 inf.sup-monoid.add-assoc*

inf.sup-monoid.add-commute)

finally have 2: $(y \sqcap z^T)^T * (-y \sqcap x) = \text{bot}$

by *simp*

have 3: *vector* $(-y)$

using *assms*(2) *vector-complement-closed* **by** *simp*

have $(-y \sqcap x)^T * (y \sqcap z^T) = (-y^T \sqcap x^T) * (y \sqcap z^T)$

by (*simp add: conv-complement conv-dist-inf*)

also have $\dots = x^T * (-y \sqcap y \sqcap z^T)$

using 3 **by** (*metis (mono-tags, opaque-lifting) conv-complement*

covector-inf-comp-3 inf.sup-monoid.add-assoc inf.sup-monoid.add-commute)

finally have 4: $(-y \sqcap x)^T * (y \sqcap z^T) = \text{bot}$

by *simp*

have 5: *univalent* $(-y \sqcap x)$

using *assms*(1) *inf-commute univalent-inf-closed* **by** *fastforce*

have $(x[y \mapsto z])^T * (x[y \mapsto z]) = (y \sqcap z^T)^T * (x[y \mapsto z]) \sqcup (-y \sqcap x)^T * (x[y \mapsto z])$

by (*simp add: conv-dist-sup mult-right-dist-sup*)

also have $\dots = (y \sqcap z^T)^T * (y \sqcap z^T) \sqcup (y \sqcap z^T)^T * (-y \sqcap x) \sqcup (-y \sqcap x)^T * (y \sqcap z^T) \sqcup (-y \sqcap x)^T * (-y \sqcap x)$

by (*simp add: mult-left-dist-sup sup-assoc*)

finally show *?thesis*

using 1 2 4 5 **by** *simp*

qed

Theorem 1.2

lemma *update-total*:

assumes *total* x

and *vector* y

and *regular* y

and *surjective* z

shows *total* $(x[y \mapsto z])$

proof –

have $(x[y \mapsto z]) * \text{top} = x * \text{top}[y \mapsto \text{top} * z]$

by (*simp add: assms*(2) *semiring.distrib-right vector-complement-closed vector-inf-comp conv-dist-comp*)

also have $\dots = \text{top}[y \mapsto \text{top}]$
using *assms(1) assms(4)* **by** *simp*
also have $\dots = \text{top}$
using *assms(3) regular-complement-top* **by** *auto*
finally show *?thesis*
by *simp*
qed

Theorem 1.3

lemma *update-mapping*:
assumes *mapping x*
and *vector y*
and *regular y*
and *bijective z*
shows *mapping (x[y ↦ z])*
using *assms update-univalent update-total* **by** *simp*

Theorem 1.4

lemma *read-injective*:
assumes *injective y*
and *univalent x*
shows *injective (x[[y]])*
using *assms injective-mult-closed univalent-conv-injective* **by** *blast*

Theorem 1.5

lemma *read-surjective*:
assumes *surjective y*
and *total x*
shows *surjective (x[[y]])*
using *assms surjective-mult-closed total-conv-surjective* **by** *blast*

Theorem 1.6

lemma *read-bijective*:
assumes *bijective y*
and *mapping x*
shows *bijective (x[[y]])*
by (*simp add: assms read-injective read-surjective*)

Theorem 1.7

lemma *read-point*:
assumes *point p*
and *mapping x*
shows *point (x[[p]])*
using *assms comp-associative read-injective read-surjective* **by** *auto*

Theorem 1.8

lemma *update-postcondition*:
assumes *point x point y*
shows $x \sqcap p = x * y^T \iff p[[x]] = y$

apply (*rule iffI*)
subgoal by (*metis assms comp-associative conv-dist-comp conv-involutive covector-inf-comp-3 equivalence-top-closed vector-covector*)
subgoal
apply (*rule order.antisym*)
subgoal by (*metis assms conv-dist-comp conv-involutive inf.boundedI inf.cobounded1 vector-covector vector-restrict-comp-conv*)
subgoal by (*smt assms comp-associative conv-dist-comp conv-involutive covector-restrict-comp-conv dense-conv-closed equivalence-top-closed inf.boundedI shunt-mapping vector-covector preorder-idempotent*)
done
done

Back and von Wright's array independence requirements [1], later also lens laws [7]

Theorem 2.1

lemma *put-get*:

assumes *vector y surjective y vector z*

shows $(x[y \multimap z])[y] = z$

proof –

have $(x[y \multimap z])[y] = (y^T \sqcap z) * y \sqcup (-y^T \sqcap x^T) * y$

by (*simp add: conv-complement conv-dist-inf conv-dist-sup mult-right-dist-sup*)

also have $\dots = z * y$

proof –

have $(-y^T \sqcap x^T) * y = \text{bot}$

by (*metis assms(1) covector-inf-comp-3 inf-commute conv-complement mult-right-zero p-inf vector-complement-closed*)

thus *?thesis*

by (*simp add: assms covector-inf-comp-3 inf-commute*)

qed

also have $\dots = z$

by (*metis assms(2,3) mult-assoc*)

finally show *?thesis*

·

qed

Theorem 2.3

lemma *put-put*:

$(x[y \multimap z])[y \multimap w] = x[y \multimap w]$

by (*metis inf-absorb2 inf-commute inf-le1 inf-sup-distrib1 maddux-3-13 sup-inf-absorb*)

Theorem 2.5

lemma *get-put*:

assumes *point y*

shows $x[y \multimap x[[y]]] = x$

proof –

have $x[y \multimap x[[y]]] = (y \sqcap y^T * x) \sqcup (-y \sqcap x)$

by (*simp add: conv-dist-comp*)

also have $\dots = (y \sqcap x) \sqcup (-y \sqcap x)$
proof –
have $y \sqcap y^T * x = y \sqcap x$
proof (*rule order.antisym*)
have $y \sqcap y^T * x = (y \sqcap y^T) * x$
by (*simp add: assms vector-inf-comp*)
also have $(y \sqcap y^T) * x = y * y^T * x$
by (*simp add: assms vector-covector*)
also have $\dots \leq x$
using *assms comp-isotone* **by** *fastforce*
finally show $y \sqcap y^T * x \leq y \sqcap x$
by *simp*
have $y \sqcap x \leq y^T * x$
by (*simp add: assms vector-restrict-comp-conv*)
thus $y \sqcap x \leq y \sqcap y^T * x$
by *simp*
qed
thus *?thesis*
by *simp*
qed
also have $\dots = x$
proof –
have *regular y*
using *assms bijective-regular* **by** *blast*
thus *?thesis*
by (*metis inf.sup-monoid.add-commute maddux-3-11-pp*)
qed
finally show *?thesis*

qed

lemma *update-inf*:

$u \leq y \implies (x[y \mapsto z]) \sqcap u = z^T \sqcap u$
by (*smt comp-inf.mult-right-dist-sup comp-inf.semiring.mult-zero-right inf.left-commute inf.sup-monoid.add-assoc inf-absorb2 p-inf sup-bot-right inf.sup-monoid.add-commute*)

lemma *update-inf-same*:

$(x[y \mapsto z]) \sqcap y = z^T \sqcap y$
by (*simp add: update-inf*)

lemma *update-inf-different*:

$u \leq -y \implies (x[y \mapsto z]) \sqcap u = x \sqcap u$
by (*smt inf.right-idem inf.sup-monoid.add-commute inf.sup-relative-same-increasing inf-import-p maddux-3-13 sup.cobounded2 update-inf-same*)

end

3 Relation-Algebraic Semantics of Disjoint-Set Forests

A disjoint-set forest represents a partition of a set into equivalence classes. We take the represented equivalence relation as the semantics of a forest. It is obtained by operation *fc* below. Additionally, operation *wcc* giving the weakly connected components of a graph will be used for the semantics of the union of two disjoint sets. Finally, operation *root* yields the root of a component tree, that is, the representative of a set containing a given element. This section defines these operations and derives their properties.

context *stone-kleene-relation-algebra*
begin

Theorem 5.2

lemma *omit-redundant-points*:

assumes *point p*

shows $p \sqcap x^* = (p \sqcap 1) \sqcup (p \sqcap x) * (-p \sqcap x)^*$

proof (*rule order.antisym*)

let $?p = p \sqcap 1$

have $?p * x * (-p \sqcap x)^* * ?p \leq ?p * top * ?p$

by (*metis comp-associative mult-left-isotone mult-right-isotone top.extremum*)

also have $... \leq ?p$

by (*simp add: assms injective-codomain vector-inf-one-comp*)

finally have $?p * x * (-p \sqcap x)^* * ?p * x \leq ?p * x$

using *mult-left-isotone* **by** *blast*

hence $?p * x * (-p \sqcap x)^* * (p \sqcap x) \leq ?p * x$

by (*simp add: assms comp-associative vector-inf-one-comp*)

also have 1: $... \leq ?p * x * (-p \sqcap x)^*$

using *mult-right-isotone star.circ-reflexive* **by** *fastforce*

finally have $?p * x * (-p \sqcap x)^* * (p \sqcap x) \sqcup ?p * x * (-p \sqcap x)^* * (-p \sqcap x) \leq ?p * x * (-p \sqcap x)^*$

by (*simp add: mult-right-isotone star.circ-plus-same star.left-plus-below-circ mult-assoc*)

hence $?p * x * (-p \sqcap x)^* * ((p \sqcup -p) \sqcap x) \leq ?p * x * (-p \sqcap x)^*$

by (*simp add: comp-inf.mult-right-dist-sup mult-left-dist-sup*)

hence $?p * x * (-p \sqcap x)^* * x \leq ?p * x * (-p \sqcap x)^*$

by (*metis assms bijective-regular inf.absorb2 inf.cobounded1 inf.sup-monoid.add-commute shunting-p*)

hence $?p * x * (-p \sqcap x)^* * x \sqcup ?p * x \leq ?p * x * (-p \sqcap x)^*$

using 1 **by** *simp*

hence $?p * (1 \sqcup x * (-p \sqcap x)^*) * x \leq ?p * x * (-p \sqcap x)^*$

by (*simp add: comp-associative mult-left-dist-sup mult-right-dist-sup*)

also have $... \leq ?p * (1 \sqcup x * (-p \sqcap x)^*)$

by (*simp add: comp-associative mult-right-isotone*)

finally have $?p * x^* \leq ?p * (1 \sqcup x * (-p \sqcap x)^*)$

using *star-right-induct* **by** (*meson dual-order.trans le-supI mult-left-sub-dist-sup-left mult-sub-right-one*)

also have $... = ?p \sqcup ?p * x * (-p \sqcap x)^*$

by (*simp add: comp-associative semiring.distrib-left*)
finally show $p \sqcap x^* \leq ?p \sqcup (p \sqcap x) * (-p \sqcap x)^*$
 by (*simp add: assms vector-inf-one-comp*)
show $?p \sqcup (p \sqcap x) * (-p \sqcap x)^* \leq p \sqcap x^*$
 by (*metis assms comp-isotone inf.boundedI inf.cobounded1 inf.coboundedI2*
inf.sup-monoid.add-commute le-supI star.circ-increasing star.circ-transitive-equal
star-isotone star-left-unfold-equal sup.cobounded1 vector-export-comp)
qed

Weakly connected components

abbreviation $wcc\ x \equiv (x \sqcup x^T)^*$

Theorem 7.1

lemma *wcc-equivalence*:
equivalence (wcc x)
apply (*intro conjI*)
subgoal by (*simp add: star.circ-reflexive*)
subgoal by (*simp add: star.circ-transitive-equal*)
subgoal by (*simp add: conv-dist-sup conv-star-commute sup-commute*)
done

Theorem 7.2

lemma *wcc-increasing*:
 $x \leq wcc\ x$
by (*simp add: star.circ-sub-dist-1*)

lemma *wcc-isotone*:
 $x \leq y \implies wcc\ x \leq wcc\ y$
using *conv-isotone star-isotone sup-mono* **by** *blast*

lemma *wcc-idempotent*:
 $wcc\ (wcc\ x) = wcc\ x$
using *star-involutive wcc-equivalence* **by** *auto*

Theorem 7.3

lemma *wcc-below-wcc*:
 $x \leq wcc\ y \implies wcc\ x \leq wcc\ y$
using *wcc-idempotent wcc-isotone* **by** *fastforce*

lemma *wcc-galois*:
 $x \leq wcc\ y \iff wcc\ x \leq wcc\ y$
using *order-trans star.circ-sub-dist-1 wcc-below-wcc* **by** *blast*

Theorem 7.4

lemma *wcc-bot*:
 $wcc\ bot = 1$
by (*simp add: star.circ-zero*)

lemma *wcc-one*:

$wcc\ 1 = 1$
by (*simp add: star-one*)

Theorem 7.5

lemma *wcc-top*:
 $wcc\ top = top$
by (*simp add: star.circ-top*)

Theorem 7.6

lemma *wcc-with-loops*:
 $wcc\ x = wcc\ (x \sqcup 1)$
by (*metis conv-dist-sup star-decompose-1 star-sup-one sup-commute symmetric-one-closed*)

lemma *wcc-without-loops*:
 $wcc\ x = wcc\ (x - 1)$
by (*metis conv-star-commute star-sum reachable-without-loops*)

lemma *forest-components-wcc*:
injective $x \implies wcc\ x = forest-components\ x$
by (*simp add: cancel-separate-1*)

Theorem 7.8

lemma *wcc-sup-wcc*:
 $wcc\ (x \sqcup y) = wcc\ (x \sqcup wcc\ y)$
by (*smt (verit, ccfv-SIG) le-sup-iff order.antisym sup-right-divisibility wcc-below-wcc wcc-increasing*)

Components of a forest, which is represented using edges directed towards the roots

abbreviation $fc\ x \equiv x^* * x^{T^*}$

Theorem 3.1

lemma *fc-equivalence*:
univalent $x \implies equivalence\ (fc\ x)$
apply (*intro conjI*)
subgoal **by** (*simp add: reflexive-mult-closed star.circ-reflexive*)
subgoal **by** (*metis cancel-separate-1 order.eq-iff star.circ-transitive-equal*)
subgoal **by** (*simp add: conv-dist-comp conv-star-commute*)
done

Theorem 3.2

lemma *fc-increasing*:
 $x \leq fc\ x$
by (*metis le-supE mult-left-isotone star.circ-back-loop-fixpoint star.circ-increasing*)

Theorem 3.3

lemma *fc-isotone*:

$x \leq y \implies fc\ x \leq fc\ y$
by (*simp add: comp-isotone conv-isotone star-isotone*)

Theorem 3.4

lemma *fc-idempotent*:

univalent $x \implies fc\ (fc\ x) = fc\ x$

by (*metis fc-equivalence cancel-separate-1 star.circ-transitive-equal star-involutive*)

Theorem 3.5

lemma *fc-star*:

univalent $x \implies (fc\ x)^* = fc\ x$

using *fc-equivalence fc-idempotent star.circ-transitive-equal* **by** *simp*

lemma *fc-plus*:

univalent $x \implies (fc\ x)^+ = fc\ x$

by (*metis fc-star star.circ-decompose-9*)

Theorem 3.6

lemma *fc-bot*:

fc bot = 1

by (*simp add: star.circ-zero*)

lemma *fc-one*:

fc 1 = 1

by (*simp add: star-one*)

Theorem 3.7

lemma *fc-top*:

fc top = top

by (*simp add: star.circ-top*)

Theorem 7.7

lemma *fc-wcc*:

univalent $x \implies wcc\ x = fc\ x$

by (*simp add: fc-star star-decompose-1*)

lemma *fc-via-root*:

assumes *total* $(p^* * (p \sqcap 1))$

shows *fc* $p = p^* * (p \sqcap 1) * p^{T^*}$

proof (*rule order.antisym*)

have $1 \leq p^* * (p \sqcap 1) * p^{T^*}$

by (*smt assms comp-associative conv-dist-comp conv-star-commute coreflexive-idempotent coreflexive-symmetric inf.cobounded2 total-var*)

hence *fc* $p \leq p^* * p^* * (p \sqcap 1) * p^{T^*} * p^{T^*}$

by (*metis comp-right-one mult-left-isotone mult-right-isotone mult-assoc*)

thus *fc* $p \leq p^* * (p \sqcap 1) * p^{T^*}$

by (*simp add: star.circ-transitive-equal mult-assoc*)

show $p^* * (p \sqcap 1) * p^{T^*} \leq fc\ p$

by (*metis comp-isotone inf.cobounded2 mult-1-right order.refl*)
qed

Theorem 5.1

lemma *update-acyclic-1*:
assumes *acyclic* ($p - 1$)
and *point* y
and *vector* w
and $w \leq p^* * y$
shows *acyclic* ($(p[w \rightarrow y]) - 1$)
proof –
let $?p = p[w \rightarrow y]$
have $w * y^T \leq p^*$
using *assms(2,4) shunt-bijective* **by** *blast*
hence $w * y^T \leq (p - 1)^*$
using *reachable-without-loops* **by** *auto*
hence $w * y^T - 1 \leq (p - 1)^* - 1$
by (*simp add: inf.coboundedI2 inf.sup-monoid.add-commute*)
also have $\dots \leq (p - 1)^+$
by (*simp add: star-plus-without-loops*)
finally have $1: w \sqcap y^T \sqcap -1 \leq (p - 1)^+$
using *assms(2,3) vector-covector* **by** *auto*
have $?p - 1 = (w \sqcap y^T \sqcap -1) \sqcup (-w \sqcap p \sqcap -1)$
by (*simp add: inf-sup-distrib2*)
also have $\dots \leq (p - 1)^+ \sqcup (-w \sqcap p \sqcap -1)$
using *1 sup-left-isotone* **by** *blast*
also have $\dots \leq (p - 1)^+ \sqcup (p - 1)$
using *comp-inf.mult-semi-associative sup-right-isotone* **by** *auto*
also have $\dots = (p - 1)^+$
by (*metis star.circ-back-loop-fixpoint sup.right-idem*)
finally have $(?p - 1)^+ \leq (p - 1)^+$
by (*metis comp-associative comp-isotone star.circ-transitive-equal star.left-plus-circ star-isotone*)
also have $\dots \leq -1$
using *assms(1)* **by** *blast*
finally show *?thesis*
by *simp*
qed

lemma *update-acyclic-2*:
assumes *acyclic* ($p - 1$)
and *point* y
and *point* x
and $y \leq p^{T*} * x$
and *univalent* p
and $p^T * y \leq y$
shows *acyclic* ($(p[p^{T*} * x \rightarrow y]) - 1$)
proof –
have $p^T * p^* * y = p^T * p * p^* * y \sqcup p^T * y$

by (*metis comp-associative mult-left-dist-sup star.circ-loop-fixpoint*)
 also have $\dots \leq p^* * y$
 by (*metis assms(5,6) comp-right-one le-supI le-supI2 mult-left-isotone star.circ-loop-fixpoint star.circ-transitive-equal*)
 finally have $p^{T^*} * x \leq p^* * y$
 by (*simp add: assms(2-4) bijective-reverse conv-star-commute comp-associative star-left-induct*)
 thus *?thesis*
 by (*simp add: assms(1-3) vector-mult-closed update-acyclic-1*)
 qed

lemma *update-acyclic-3*:
 assumes *acyclic* ($p - 1$)
 and *point* y
 and *point* w
 and $y \leq p^{T^*} * w$
 shows *acyclic* ($(p[w \longrightarrow y]) - 1$)
 by (*simp add: assms bijective-reverse conv-star-commute update-acyclic-1*)

lemma *rectangle-star-rectangle*:
rectangle $a \implies a * x^* * a \leq a$
 by (*metis mult-left-isotone mult-right-isotone top.extremum*)

lemma *arc-star-arc*:
arc $a \implies a * x^* * a \leq a$
 using *arc-top-arc rectangle-star-rectangle* by *blast*

lemma *star-rectangle-decompose*:
 assumes *rectangle* a
 shows $(a \sqcup x)^* = x^* \sqcup x^* * a * x^*$
proof (*rule order.antisym*)
 have $1: 1 \leq x^* \sqcup x^* * a * x^*$
 by (*simp add: star.circ-reflexive sup.coboundedI1*)
 have $(a \sqcup x) * (x^* \sqcup x^* * a * x^*) = a * x^* \sqcup a * x^* * a * x^* \sqcup x^+ \sqcup x^+ * a * x^*$
 by (*metis comp-associative semiring.combine-common-factor semiring.distrib-left sup-commute*)
 also have $\dots = a * x^* \sqcup x^+ \sqcup x^+ * a * x^*$
 using *assms rectangle-star-rectangle* by (*simp add: mult-left-isotone sup-absorb1*)
 also have $\dots = x^+ \sqcup x^* * a * x^*$
 by (*metis comp-associative star.circ-loop-fixpoint sup-assoc sup-commute*)
 also have $\dots \leq x^* \sqcup x^* * a * x^*$
 using *star.left-plus-below-circ sup-left-isotone* by *auto*
 finally show $(a \sqcup x)^* \leq x^* \sqcup x^* * a * x^*$
 using 1 by (*metis comp-right-one le-supI star-left-induct*)
next
 show $x^* \sqcup x^* * a * x^* \leq (a \sqcup x)^*$
 by (*metis comp-isotone le-supE le-supI star.circ-increasing*)

star.circ-transitive-equal star-isotone sup-ge2)
qed

lemma *star-arc-decompose*:

arc a $\implies (a \sqcup x)^* = x^* \sqcup x^* * a * x^*$
using *arc-top-arc star-rectangle-decompose* **by** *blast*

lemma *plus-rectangle-decompose*:

assumes *rectangle a*
shows $(a \sqcup x)^+ = x^+ \sqcup x^* * a * x^*$

proof –

have $(a \sqcup x)^+ = (a \sqcup x) * (x^* \sqcup x^* * a * x^*)$
by (*simp add: assms star-rectangle-decompose*)
also have $\dots = a * x^* \sqcup a * x^* * a * x^* \sqcup x^+ \sqcup x^+ * a * x^*$
by (*metis comp-associative semiring.combine-common-factor*

semiring.distrib-left sup-commute)

also have $\dots = a * x^* \sqcup x^+ \sqcup x^+ * a * x^*$
using *assms rectangle-star-rectangle* **by** (*simp add: mult-left-isotone*
sup-absorb1)

also have $\dots = x^+ \sqcup x^* * a * x^*$
by (*metis comp-associative star.circ-loop-fixpoint sup-assoc sup-commute*)

finally show *?thesis*

by *simp*

qed

Theorem 8.1

lemma *plus-arc-decompose*:

arc a $\implies (a \sqcup x)^+ = x^+ \sqcup x^* * a * x^*$
using *arc-top-arc plus-rectangle-decompose* **by** *blast*

Theorem 8.2

lemma *update-acyclic-4*:

assumes *acyclic (p - 1)*
and *point y*
and *point w*
and $y \sqcap p^* * w = \text{bot}$
shows *acyclic ((p[w \longrightarrow y]) - 1)*

proof –

let $?p = p[w \longrightarrow y]$
have $y^T * p^* * w \leq -1$
using *assms(4) comp-associative pseudo-complement schroeder-3-p* **by** *auto*
hence $1: p^* * w * y^T * p^* \leq -1$
by (*metis comp-associative comp-commute-below-diversity*

star.circ-transitive-equal)

have $?p - 1 \leq (w \sqcap y^T) \sqcup (p - 1)$

by (*metis comp-inf.mult-right-dist-sup dual-order.trans inf.cobounded1*
inf.coboundedI2 inf.sup-monoid.add-assoc le-supI sup.cobounded1 sup-ge2)

also have $\dots = w * y^T \sqcup (p - 1)$

using *assms(2,3)* **by** (*simp add: vector-covector*)

finally have $(?p - 1)^+ \leq (w * y^T \sqcup (p - 1))^+$
by (*simp add: comp-isotone star-isotone*)
also have $\dots = (p - 1)^+ \sqcup (p - 1)^* * w * y^T * (p - 1)^*$
using *assms(2,3) plus-arc-decompose points-arc* **by** (*simp add: comp-associative*)
also have $\dots \leq (p - 1)^+ \sqcup p^* * w * y^T * p^*$
using *reachable-without-loops* **by** *auto*
also have $\dots \leq -1$
using *1 assms(1)* **by** *simp*
finally show *?thesis*
by *simp*
qed

Theorem 8.3

lemma *update-acyclic-5*:
assumes *acyclic (p - 1)*
and *point w*
shows *acyclic ((p[w→w]) - 1)*
proof –
let $?p = p[w \rightarrow w]$
have $?p - 1 \leq (w \sqcap w^T \sqcap -1) \sqcup (p - 1)$
by (*metis comp-inf.mult-right-dist-sup inf.cobounded2 inf.sup-monoid.add-assoc sup-right-isotone*)
also have $\dots = p - 1$
using *assms(2)* **by** (*metis comp-inf.covector-complement-closed equivalence-top-closed inf-top.right-neutral maddux-3-13 pseudo-complement regular-closed-top regular-one-closed vector-covector vector-top-closed*)
finally show *?thesis*
using *assms(1) acyclic-down-closed* **by** *blast*
qed

Root of the tree containing point x in the disjoint-set forest p

abbreviation $roots\ p \equiv (p \sqcap 1) * top$
abbreviation $root\ p\ x \equiv p^{T*} * x \sqcap roots\ p$

Theorem 4.1

lemma *root-var*:
 $root\ p\ x = (p \sqcap 1) * p^{T*} * x$
by (*simp add: coreflexive-comp-top-inf inf-commute mult-assoc*)

Theorem 4.2

lemma *root-successor-loop*:
 $univalent\ p \implies root\ p\ x = p[[root\ p\ x]]$
by (*metis root-var injective-codomain comp-associative conv-dist-inf coreflexive-symmetric equivalence-one-closed inf.cobounded2 univalent-conv-injective*)

lemma *root-transitive-successor-loop*:
 $univalent\ p \implies root\ p\ x = p^{T*} * (root\ p\ x)$

by (*metis mult-1-right star-one star-simulation-right-equal root-successor-loop*)

lemma *roots-successor-loop*:

univalent p $\implies p[[\text{roots } p]] = \text{roots } p$

by (*metis conv-involutive inf-commute injective-codomain one-inf-conv mult-assoc*)

lemma *roots-transitive-successor-loop*:

univalent p $\implies p^{T^*} * (\text{roots } p) = \text{roots } p$

by (*metis comp-associative star.circ-left-top star-simulation-right-equal roots-successor-loop*)

The root of a tree of a node belongs to the same component as the node.

lemma *root-same-component*:

injective x $\implies \text{root } p \ x * x^T \leq \text{fc } p$

by (*metis comp-associative coreflexive-comp-top-inf eq-refl inf.sup-left-divisibility inf.sup-monoid.add-commute mult-isotone star.circ-circ-mult star.circ-right-top star.circ-transitive-equal star-one star-outer-increasing test-preserves-equation top-greatest*)

lemma *root-vector*:

vector x $\implies \text{vector } (\text{root } p \ x)$

by (*simp add: vector-mult-closed root-var*)

lemma *root-vector-inf*:

vector x $\implies \text{root } p \ x * x^T = \text{root } p \ x \sqcap x^T$

by (*simp add: vector-covector root-vector*)

lemma *root-same-component-vector*:

injective x $\implies \text{vector } x \implies \text{root } p \ x \sqcap x^T \leq \text{fc } p$

using *root-same-component root-vector-inf* **by** *fastforce*

lemma *univalent-root-successors*:

assumes *univalent p*

shows $(p \sqcap 1) * p^* = p \sqcap 1$

proof (*rule order.antisym*)

have $(p \sqcap 1) * p \leq p \sqcap 1$

by (*smt assms(1) comp-inf.mult-semi-associative conv-dist-comp conv-dist-inf conv-order equivalence-one-closed inf.absorb1 inf.sup-monoid.add-assoc injective-codomain*)

thus $(p \sqcap 1) * p^* \leq p \sqcap 1$

using *star-right-induct-mult* **by** *blast*

show $p \sqcap 1 \leq (p \sqcap 1) * p^*$

by (*metis coreflexive-idempotent inf-le1 inf-le2 mult-right-isotone order-trans star.circ-increasing*)

qed

lemma *same-component-same-root-sub*:

assumes *univalent p*

and *bijjective* y
and $x * y^T \leq_{fc} p$
shows $root\ p\ x \leq root\ p\ y$
proof –
have $root\ p\ x * y^T \leq (p \sqcap 1) * p^{T*}$
by (*smt* *assms*(1,3) *mult-isotone* *mult-assoc* *root-var* *fc-plus* *fc-star* *order.eq-iff*
univalent-root-successors)
thus *?thesis*
by (*simp* *add: assms*(2) *shunt-bijjective* *root-var*)
qed

lemma *same-component-same-root:*

assumes *univalent* p
and *bijjective* x
and *bijjective* y
and $x * y^T \leq_{fc} p$
shows $root\ p\ x = root\ p\ y$
proof (*rule* *order.antisym*)
show $root\ p\ x \leq root\ p\ y$
using *assms*(1,3,4) *same-component-same-root-sub* **by** *blast*
have $y * x^T \leq_{fc} p$
using *assms*(1,4) *fc-equivalence* *conv-dist-comp* *conv-isotone* **by** *fastforce*
thus $root\ p\ y \leq root\ p\ x$
using *assms*(1,2) *same-component-same-root-sub* **by** *blast*
qed

lemma *same-roots-sub:*

assumes *univalent* q
and $p \sqcap 1 \leq q \sqcap 1$
and $fc\ p \leq_{fc} q$
shows $p^* * (p \sqcap 1) \leq q^* * (q \sqcap 1)$
proof –
have $p^* * (p \sqcap 1) \leq p^* * (q \sqcap 1)$
using *assms*(2) *mult-right-isotone* **by** *auto*
also have $\dots \leq_{fc} p * (q \sqcap 1)$
using *mult-left-isotone* *mult-right-isotone* *star.circ-reflexive* **by** *fastforce*
also have $\dots \leq_{fc} q * (q \sqcap 1)$
by (*simp* *add: assms*(3) *mult-left-isotone*)
also have $\dots = q^* * (q \sqcap 1)$
by (*metis* *assms*(1) *conv-dist-comp* *conv-dist-inf* *conv-star-commute*
inf-commute *one-inf-conv* *symmetric-one-closed* *mult-assoc*
univalent-root-successors)
finally show *?thesis*

qed

lemma *same-roots:*

assumes *univalent* p
and *univalent* q

and $p \sqcap 1 = q \sqcap 1$
and $fc\ p = fc\ q$
shows $p^* * (p \sqcap 1) = q^* * (q \sqcap 1)$
by (*smt assms conv-dist-comp conv-dist-inf conv-involutive conv-star-commute inf-commute one-inf-conv symmetric-one-closed root-var univalent-root-successors*)

lemma *same-root*:

assumes *univalent p*
and *univalent q*
and $p \sqcap 1 = q \sqcap 1$
and $fc\ p = fc\ q$
shows $root\ p\ x = root\ q\ x$
by (*metis assms mult-assoc root-var univalent-root-successors*)

lemma *loop-root*:

assumes *injective x*
and $x = p[[x]]$
shows $x = root\ p\ x$
proof (*rule order.antisym*)
have $x \leq p * x$
by (*metis assms comp-associative comp-right-one conv-order equivalence-one-closed ex231c inf.orderE inf.sup-monoid.add-commute mult-left-isotone mult-right-isotone one-inf-conv*)
hence $x = (p \sqcap 1) * x$
by (*simp add: assms(1) inf-absorb2 injective-comp-right-dist-inf*)
thus $x \leq root\ p\ x$
by (*metis assms(2) coreflexive-comp-top-inf inf.boundedI inf.cobounded1 inf.cobounded2 mult-isotone star.circ-increasing*)
next
show $root\ p\ x \leq x$
using *assms(2) le-infI1 star-left-induct-mult* **by** *auto*
qed

lemma *one-loop*:

assumes *acyclic (p - 1)*
and *univalent p*
shows $(p \sqcap 1) * (p^T - 1)^+ * (p \sqcap 1) = bot$
proof -
have $p^{T+} \sqcap (p \sqcap 1) * top * (p \sqcap 1) = (p \sqcap 1) * p^{T+} * (p \sqcap 1)$
by (*simp add: test-comp-test-top*)
also have $\dots \leq p^{T*} * (p \sqcap 1)$
by (*simp add: inf.coboundedI2 mult-left-isotone star.circ-mult-upper-bound star.circ-reflexive star.left-plus-below-circ*)
also have $\dots = p \sqcap 1$
by (*metis assms(2) conv-dist-comp conv-dist-inf conv-star-commute inf-commute one-inf-conv symmetric-one-closed univalent-root-successors*)
also have $\dots \leq 1$
by *simp*
finally have $(p \sqcap 1) * top * (p \sqcap 1) \leq -(p^{T+} - 1)$

using *p-antitone p-antitone-iff p-shunting-swap* **by** *blast*
hence $(p \sqcap 1)^T * (p^{T+} - 1) * (p \sqcap 1)^T \leq \text{bot}$
using *triple-schroeder-p p-top* **by** *blast*
hence $(p \sqcap 1) * (p^{T+} - 1) * (p \sqcap 1) = \text{bot}$
by (*simp add: coreflexive-symmetric le-bot*)
thus *?thesis*
by (*smt assms(1) conv-complement conv-dist-comp conv-dist-inf*
conv-star-commute inf-absorb1 star.circ-plus-same symmetric-one-closed
reachable-without-loops star-plus-without-loops)
qed

lemma *root-root*:

$\text{root } p \ x = \text{root } p \ (\text{root } p \ x)$
by (*smt comp-associative comp-inf.mult-right-sub-dist-sup-right dual-order.eq-iff*
inf.cobounded1 inf.cobounded2 inf.orderE mult-right-isotone
star.circ-loop-fixpoint star.circ-transitive-equal root-var)

lemma *loop-root-2*:

assumes *acyclic (p - 1)*
and *univalent p*
and *injective x*
and $x \leq p^{T+} * x$
shows $x = \text{root } p \ x$
proof (*rule order.antisym*)
have $1: x = x - (-1 * x)$
by (*metis assms(3) comp-injective-below-complement inf.orderE mult-1-left*
regular-one-closed)
have $x \leq (p^T - 1)^+ * x \sqcup (p \sqcap 1) * x$
by (*metis assms(4) inf-commute mult-right-dist-sup one-inf-conv*
plus-reachable-without-loops)
also have $\dots \leq -1 * x \sqcup (p \sqcap 1) * x$
by (*metis assms(1) conv-complement conv-dist-inf conv-isotone*
conv-plus-commute mult-left-isotone semiring.add-right-mono
symmetric-one-closed)
also have $\dots \leq -1 * x \sqcup \text{root } p \ x$
using *comp-isotone inf.coboundedI2 star.circ-reflexive sup-right-isotone* **by**
auto
finally have $x \leq (-1 * x \sqcup \text{root } p \ x) - (-1 * x)$
using 1 *inf.boundedI inf.order-iff* **by** *blast*
also have $\dots \leq \text{root } p \ x$
using *inf.sup-left-divisibility* **by** *auto*
finally show $2: x \leq \text{root } p \ x$
 \cdot
have $\text{root } p \ x = (p \sqcap 1) * x \sqcup (p \sqcap 1) * (p^T - 1)^+ * x$
by (*metis comp-associative mult-left-dist-sup star.circ-loop-fixpoint*
sup-commute reachable-without-loops root-var)
also have $\dots \leq x \sqcup (p \sqcap 1) * (p^T - 1)^+ * \text{root } p \ x$
using 2 **by** (*metis coreflexive-comp-top-inf inf.cobounded2 mult-right-isotone*
semiring.add-mono)

also have $\dots = x$
by (*metis assms(1,2) one-loop root-var mult-assoc semiring.mult-not-zero sup-bot-right*)
finally show $\text{root } p \ x \leq x$
qed

lemma *path-compression-invariant-simplify*:

assumes *point w*
and $p^{T+} * w \leq -w$
and $w \neq y$
shows $p[[w]] \neq w$
proof
assume $p[[w]] = w$
hence $w \leq p^{T+} * w$
by (*metis comp-isotone eq-refl star.circ-mult-increasing*)
also have $\dots \leq -w$
by (*simp add: assms(2)*)
finally have $w = \text{bot}$
using *inf.orderE* **by** *fastforce*
thus *False*
using *assms(1,3) le-bot* **by** *force*
qed

end

context *stone-relation-algebra-tarski*
begin

Theorem 5.4 *distinct-points* has been moved to theory *Relation-Algebras* in entry *Stone-Relation-Algebras*

[Back and von Wright's array independence requirements \[1\]](#)

[Theorem 2.2](#)

lemma *put-get-different-vector*:

assumes *vector y w ≤ -y*
shows $(x[y \mapsto z])[[w]] = x[[w]]$
proof –
have $(x[y \mapsto z])[[w]] = (y^T \sqcap z) * w \sqcup (-y^T \sqcap x^T) * w$
by (*simp add: conv-complement conv-dist-inf conv-dist-sup mult-right-dist-sup*)
also have $\dots = z * (w \sqcap y) \sqcup x^T * (w - y)$
by (*metis assms(1) conv-complement covector-inf-comp-3 inf-commute vector-complement-closed*)
also have $\dots = z * (w \sqcap y) \sqcup x^T * w$
by (*simp add: assms(2) inf.absorb1*)
also have $\dots = z * \text{bot} \sqcup x^T * w$
by (*metis assms(2) comp-inf.semiring.mult-zero-right inf.absorb1 inf.sup-monoid.add-assoc p-inf*)
also have $\dots = x^T * w$

by *simp*
 finally show *?thesis*
 .
 qed

lemma *put-get-different*:
 assumes *point y point w w ≠ y*
 shows $(x[y \mapsto z])[w] = x[w]$
proof –
 have $w \sqcap y = \text{bot}$
 using *assms distinct-points by simp*
 hence $w \leq -y$
 using *pseudo-complement by simp*
 thus *?thesis*
 by (*simp add: assms(1) assms(2) put-get-different-vector*)
 qed

Theorem 2.4

lemma *put-put-different-vector*:
 assumes *vector y vector v v ⊓ y = bot*
 shows $(x[y \mapsto z])[v \mapsto w] = (x[v \mapsto w])[y \mapsto z]$
proof –
 have $(x[y \mapsto z])[v \mapsto w] = (v \sqcap w^T) \sqcup (-v \sqcap y \sqcap z^T) \sqcup (-v \sqcap -y \sqcap x)$
 by (*simp add: comp-inf.semiring.distrib-left inf-assoc sup-assoc*)
 also have $\dots = (v \sqcap w^T) \sqcup (y \sqcap z^T) \sqcup (-v \sqcap -y \sqcap x)$
 by (*metis assms(3) inf-commute inf-import-p p-inf selection-closed-id*)
 also have $\dots = (y \sqcap z^T) \sqcup (v \sqcap w^T) \sqcup (-y \sqcap -v \sqcap x)$
 by (*simp add: inf-commute sup-commute*)
 also have $\dots = (y \sqcap z^T) \sqcup (-y \sqcap v \sqcap w^T) \sqcup (-y \sqcap -v \sqcap x)$
 using *assms distinct-points pseudo-complement inf.absorb2 by simp*
 also have $\dots = (x[v \mapsto w])[y \mapsto z]$
 by (*simp add: comp-inf.semiring.distrib-left inf-assoc sup-assoc*)
 finally show *?thesis*
 .
 qed

lemma *put-put-different*:
 assumes *point y point v v ≠ y*
 shows $(x[y \mapsto z])[v \mapsto w] = (x[v \mapsto w])[y \mapsto z]$
 using *assms distinct-points put-put-different-vector by blast*

end

4 Verifying Operations on Disjoint-Set Forests

In this section we verify the make-set, find-set and union-sets operations of disjoint-set forests. We start by introducing syntax for updating arrays in programs. Updating the value at a given array index means updating the

whole array.

syntax

-rel-update :: *idt* \Rightarrow 'a \Rightarrow 'a \Rightarrow 'b *com* ((λ [-] :=/ -) [70, 65, 65] 61)

translations

$x[y] := z \Rightarrow (x := (y \sqcap z^T) \sqcup (CONST \text{uminus } y \sqcap x))$

The finiteness requirement in the following class is used for proving that the operations terminate.

class *finite-regular-p-algebra* = *p-algebra* +
assumes *finite-regular*: *finite* { *x* . *regular* *x* }
begin

abbreviation *card-down-regular* :: 'a \Rightarrow *nat* (\downarrow [100] 100)
where $x \downarrow \equiv \text{card } \{ z . \text{regular } z \wedge z \leq x \}$

end

class *stone-kleene-relation-algebra-tarski-finite-regular* =
stone-kleene-relation-algebra-tarski + *finite-regular-p-algebra*
begin

4.1 Make-Set

We prove two correctness results about make-set. The first shows that the forest changes only to the extent of making one node the root of a tree. The second result adds that only singleton sets are created.

definition *make-set-postcondition* $p \ x \ p0 \equiv x \sqcap p = x * x^T \wedge -x \sqcap p = -x \sqcap p0$

theorem *make-set*:

VARs *p*
[*point* $x \wedge p0 = p$]
 $p[x] := x$
[*make-set-postcondition* $p \ x \ p0$]
apply *vcg-tc-simp*
by (*simp* *add*: *make-set-postcondition-def* *inf-sup-distrib1* *inf-assoc*[*THEN sym*]
vector-covector[*THEN sym*])

theorem *make-set-2*:

VARs *p*
[*point* $x \wedge p0 = p \wedge p \leq 1$]
 $p[x] := x$
[*make-set-postcondition* $p \ x \ p0 \wedge p \leq 1$]

proof *vcg-tc*

fix *p*

assume *1*: *point* $x \wedge p0 = p \wedge p \leq 1$

show *make-set-postcondition* ($p[x \mapsto x]$) $x \ p0 \wedge p[x \mapsto x] \leq 1$

proof (*rule conjI*)

```

show make-set-postcondition (p[x $\rightarrow$ x]) x p0
  using 1 by (simp add: make-set-postcondition-def inf-sup-distrib1
inf-assoc[THEN sym] vector-covector[THEN sym])
  show p[x $\rightarrow$ x]  $\leq$  1
  using 1 by (metis coreflexive-sup-closed dual-order.trans inf.cobounded2
vector-covector)
qed
qed

```

The above total-correctness proof allows us to extract a function, which can be used in other implementations below. This is a technique of [10].

```

lemma make-set-exists:
  point x  $\implies$   $\exists$  p' . make-set-postcondition p' x p
  using tc-extract-function make-set by blast

```

```

definition make-set p x  $\equiv$  (SOME p' . make-set-postcondition p' x p)

```

```

lemma make-set-function:
  assumes point x
  and p' = make-set p x
  shows make-set-postcondition p' x p
proof –
  let ?P =  $\lambda$ p' . make-set-postcondition p' x p
  have ?P (SOME z . ?P z)
  using assms(1) make-set-exists by (meson someI)
  thus ?thesis
  using assms(2) make-set-def by auto
qed

```

end

4.2 Find-Set

Disjoint-set forests are represented by their parent mapping. It is a forest except each root of a component tree points to itself.

We prove that `find-set` returns the root of the component tree of the given node.

```

context pd-kleene-allegory
begin

```

```

abbreviation disjoint-set-forest p  $\equiv$  mapping p  $\wedge$  acyclic (p – 1)

```

end

```

context stone-kleene-relation-algebra-tarski-finite-regular
begin

```

```

definition find-set-precondition p x  $\equiv$  disjoint-set-forest p  $\wedge$  point x

```

definition *find-set-invariant* $p x y \equiv \text{find-set-precondition } p x \wedge \text{point } y \wedge y \leq p^{T^*} * x$

definition *find-set-postcondition* $p x y \equiv \text{point } y \wedge y = \text{root } p x$

lemma *find-set-1*:

find-set-precondition $p x \implies \text{find-set-invariant } p x x$

apply (*unfold find-set-invariant-def*)

using *mult-left-isotone star.circ-reflexive find-set-precondition-def* **by** *fastforce*

lemma *find-set-2*:

find-set-invariant $p x y \wedge y \neq p[[y]] \implies \text{find-set-invariant } p x (p[[y]]) \wedge (p^{T^*} * (p[[y]])) \downarrow < (p^{T^*} * y) \downarrow$

proof –

let $?s = \{ z . \text{regular } z \wedge z \leq p^{T^*} * y \}$

let $?t = \{ z . \text{regular } z \wedge z \leq p^{T^*} * (p[[y]]) \}$

assume 1: *find-set-invariant* $p x y \wedge y \neq p[[y]]$

have 2: *point* $(p[[y]])$

using 1 *read-point find-set-invariant-def find-set-precondition-def* **by** *simp*

show *find-set-invariant* $p x (p[[y]]) \wedge \text{card } ?t < \text{card } ?s$

proof (*unfold find-set-invariant-def, intro conjI*)

show *find-set-precondition* $p x$

using 1 *find-set-invariant-def* **by** *simp*

show *vector* $(p[[y]])$

using 2 **by** *simp*

show *injective* $(p[[y]])$

using 2 **by** *simp*

show *surjective* $(p[[y]])$

using 2 **by** *simp*

show $p[[y]] \leq p^{T^*} * x$

using 1 **by** (*metis (opaque-lifting) find-set-invariant-def comp-associative comp-isotone star.circ-increasing star.circ-transitive-equal*)

show $\text{card } ?t < \text{card } ?s$

proof –

have $p[[y]] = (p^T \sqcap 1) * y \sqcup (p^T - 1) * y$

by (*metis maddux-3-11-pp mult-right-dist-sup regular-one-closed*)

also have $\dots \leq ((p[[y]]) \sqcap y) \sqcup (p^T - 1) * y$

by (*metis comp-left-subdist-inf mult-1-left semiring.add-right-mono*)

also have $\dots = (p^T - 1) * y$

using 1 2 *find-set-invariant-def distinct-points* **by** *auto*

finally have 3: $(p^T - 1)^* * (p[[y]]) \leq (p^T - 1)^+ * y$

by (*simp add: mult-right-isotone star-simulation-right-equal mult-assoc*)

have $p^{T^*} * (p[[y]]) \leq p^{T^*} * y$

by (*metis mult-left-isotone star.right-plus-below-circ mult-assoc*)

hence 4: $?t \subseteq ?s$

using *order-trans* **by** *auto*

have 5: $y \in ?s$

using 1 *find-set-invariant-def bijective-regular mult-left-isotone*

star.circ-reflexive **by** *fastforce*

have 6: $\neg y \in ?t$


```

proof
  assume  $y \in ?t$ 
  hence  $y \leq (p^T - 1)^+ * y$ 
    using  $\mathcal{J}$  by (metis reachable-without-loops mem-Collect-eq order-trans)
  hence  $y * y^T \leq (p^T - 1)^+$ 
    using 1 find-set-invariant-def shunt-bijective by simp
  also have  $\dots \leq -1$ 
    using 1 by (metis (mono-tags, lifting) find-set-invariant-def
find-set-precondition-def conv-dist-comp conv-dist-inf conv-isotone
conv-star-commute equivalence-one-closed star.circ-plus-same
symmetric-complement-closed)
  finally have  $y \leq -y$ 
    using schroeder-4-p by auto
  thus False
    using 1 by (metis find-set-invariant-def comp-inf.coreflexive-idempotent
conv-complement covector-vector-comp inf.absorb1 inf.sup-monoid.add-commute
pseudo-complement surjective-conv-total top.extremum vector-top-closed
regular-closed-top)
  qed
  show  $\text{card } ?t < \text{card } ?s$ 
    apply (rule psubset-card-mono)
    subgoal using finite-regular by simp
    subgoal using 4 5 6 by auto
    done
  qed
qed
qed

```

lemma *find-set-3*:

```

find-set-invariant  $p \ x \ y \wedge y = p[[y]] \implies \text{find-set-postcondition } p \ x \ y$ 
proof –
  assume 1: find-set-invariant  $p \ x \ y \wedge y = p[[y]]$ 
  show find-set-postcondition  $p \ x \ y$ 
  proof (unfold find-set-postcondition-def, rule conjI)
    show point  $y$ 
      using 1 find-set-invariant-def by simp
    show  $y = \text{root } p \ x$ 
    proof (rule order.antisym)
      have  $y * y^T \leq p$ 
        using 1 by (metis find-set-invariant-def find-set-precondition-def
shunt-bijective shunt-mapping top-right-mult-increasing)
      hence  $y * y^T \leq p \sqcap 1$ 
        using 1 find-set-invariant-def le-infI by blast
      hence  $y \leq \text{roots } p$ 
        using 1 by (metis find-set-invariant-def order-lesseq-imp shunt-bijective
top-right-mult-increasing mult-assoc)
      thus  $y \leq \text{root } p \ x$ 
        using 1 find-set-invariant-def by simp
    next

```

```

have 2:  $x \leq p^* * y$ 
  using 1 find-set-invariant-def find-set-precondition-def bijective-reverse
conv-star-commute by auto
  have  $p^T * p^* * y = p^T * p * p^* * y \sqcup (p[[y]])$ 
    by (metis comp-associative mult-left-dist-sup star.circ-loop-fixpoint)
  also have  $\dots \leq p^* * y \sqcup y$ 
    using 1 by (metis find-set-invariant-def find-set-precondition-def
comp-isotone mult-left-sub-dist-sup semiring.add-right-mono
star.circ-back-loop-fixpoint star.circ-circ-mult star.circ-top
star.circ-transitive-equal star-involutive star-one)
  also have  $\dots = p^* * y$ 
    by (metis star.circ-loop-fixpoint sup.left-idem sup-commute)
  finally have 3:  $p^{T*} * x \leq p^* * y$ 
    using 2 by (simp add: comp-associative star-left-induct)
  have  $p * y \sqcap \text{roots } p = (p \sqcap 1) * p * y$ 
    using comp-associative coreflexive-comp-top-inf inf-commute by auto
  also have  $\dots \leq p^T * p * y$ 
    by (metis inf.cobounded2 inf.sup-monoid.add-commute mult-left-isotone
one-inf-conv)
  also have  $\dots \leq y$ 
    using 1 find-set-invariant-def find-set-precondition-def mult-left-isotone by
fastforce
  finally have 4:  $p * y \leq y \sqcup -\text{roots } p$ 
    using 1 by (metis find-set-invariant-def shunting-p bijective-regular)
  have  $p * -\text{roots } p \leq -\text{roots } p$ 
    using 1 by (metis find-set-invariant-def find-set-precondition-def
conv-complement-sub-leq conv-involutive roots-successor-loop)
  hence  $p * y \sqcup p * -\text{roots } p \leq y \sqcup -\text{roots } p$ 
    using 4 dual-order.trans le-supI sup-ge2 by blast
  hence  $p * (y \sqcup -\text{roots } p) \leq y \sqcup -\text{roots } p$ 
    by (simp add: mult-left-dist-sup)
  hence  $p^* * y \leq y \sqcup -\text{roots } p$ 
    by (simp add: star-left-induct)
  hence  $p^{T*} * x \leq y \sqcup -\text{roots } p$ 
    using 3 dual-order.trans by blast
  thus  $\text{root } p \ x \leq y$ 
    using 1 by (metis find-set-invariant-def shunting-p bijective-regular)
qed
qed
qed

```

```

theorem find-set:
  VARs  $y$ 
  [ find-set-precondition  $p \ x$  ]
   $y := x;$ 
  WHILE  $y \neq p[[y]]$ 
  INV { find-set-invariant  $p \ x \ y$  }
  VAR {  $(p^{T*} * y) \downarrow$  }
  DO  $y := p[[y]]$ 

```

```

    OD
  [ find-set-postcondition p x y ]
  apply vcg-tc-simp
    apply (fact find-set-1)
    apply (fact find-set-2)
  by (fact find-set-3)

```

lemma *find-set-exists*:
 $find\text{-set}\text{-precondition } p \ x \implies \exists y . find\text{-set}\text{-postcondition } p \ x \ y$
using *tc-extract-function find-set by blast*

The root of a component tree is a point, that is, represents a singleton set of nodes. This could be proved from the definitions using Kleene-relation algebraic calculations. But they can be avoided because the property directly follows from the postcondition of the previous correctness proof. The corresponding algorithm shows how to obtain the root. We therefore have an essentially constructive proof of the following result.

Theorem 4.3

lemma *root-point*:
 $disjoint\text{-set}\text{-forest } p \implies point \ x \implies point \ (root \ p \ x)$
using *find-set-exists find-set-precondition-def find-set-postcondition-def by simp*

definition *find-set* $p \ x \equiv (SOME \ y . find\text{-set}\text{-postcondition } p \ x \ y)$

lemma *find-set-function*:
assumes *find-set-precondition* $p \ x$
and $y = find\text{-set} \ p \ x$
shows *find-set-postcondition* $p \ x \ y$
by (*metis assms find-set-def find-set-exists someI*)

4.3 Path Compression

The path-compression technique is frequently implemented in recursive implementations of find-set modifying the tree on the way out from recursive calls. Here we implement it using a second while-loop, which iterates over the same path to the root and changes edges to point to the root of the component, which is known after the while-loop in find-set completes. We prove that path compression preserves the equivalence-relational semantics of the disjoint-set forest and also preserves the roots of the component trees. Additionally we prove the exact effect of path compression.

definition *path-compression-precondition* $p \ x \ y \equiv disjoint\text{-set}\text{-forest } p \ \wedge \ point \ x \ \wedge \ point \ y \ \wedge \ y = root \ p \ x$

definition *path-compression-invariant* $p \ x \ y \ p0 \ w \equiv$
 $path\text{-compression}\text{-precondition } p \ x \ y \ \wedge \ point \ w \ \wedge \ y \leq p^{T^*} * w \ \wedge$
 $(w \neq x \longrightarrow p[[x]] = y \ \wedge \ y \neq x \ \wedge \ p^{T^+} * w \leq -x) \ \wedge \ p \sqcap 1 = p0 \sqcap 1 \ \wedge \ fc \ p =$
 $fc \ p0 \ \wedge$

$\text{root } p \ w = y \wedge (w \neq y \longrightarrow p^{T^+} * w \leq -w) \wedge p[[w]] = p0[[w]] \wedge p0[p0^{T^*} * x$
 $- p0^{T^*} * w \longrightarrow y] = p \wedge$
 $\text{disjoint-set-forest } p0 \wedge y = \text{root } p0 \ x \wedge w \leq p0^{T^*} * x$
definition *path-compression-postcondition* $p \ x \ y \ p0 \equiv$
path-compression-precondition $p \ x \ y \wedge p \sqcap 1 = p0 \sqcap 1 \wedge \text{fc } p = \text{fc } p0 \wedge$
 $p0[p0^{T^*} * x \longrightarrow y] = p$

We first consider a variant that achieves the effect as a single update. The parents of all nodes reachable from x are simultaneously updated to the root of the component of x .

lemma *path-compression-exact*:

assumes *path-compression-precondition* $p0 \ x \ y$

and $p0[p0^{T^*} * x \longrightarrow y] = p$

shows $p \sqcap 1 = p0 \sqcap 1 \ \text{fc } p = \text{fc } p0$

proof –

have $a1$: *disjoint-set-forest* $p0$ **and** $a2$: *point* x **and** $a3$: *point* y **and** $a4$: $y = \text{root } p0 \ x$

using *path-compression-precondition-def* *assms(1)* **by** *auto*

have 1 : *regular* $(p0^{T^*} * x)$

using $a1 \ a2$ *bijjective-regular mapping-regular regular-closed-star regular-conv-closed regular-mult-closed* **by** *auto*

have $p \sqcap 1 = (p0^{T^*} * x \sqcap y^T \sqcap 1) \sqcup (-(p0^{T^*} * x) \sqcap p0 \sqcap 1)$

using *assms(2)* *inf-sup-distrib2* **by** *auto*

also have $\dots = (p0^{T^*} * x \sqcap p0 \sqcap 1) \sqcup (-(p0^{T^*} * x) \sqcap p0 \sqcap 1)$

proof –

have $p0^{T^*} * x \sqcap y^T \sqcap 1 = p0^{T^*} * x \sqcap p0 \sqcap 1$

proof (*rule order.antisym*)

have $(p0 \sqcap 1) * p0^{T^*} * x \sqcap 1 \leq p0$

by (*smt coreflexive-comp-top-inf-one inf.absorb-iff2 inf.cobounded2 inf.sup-monoid.add-assoc root-var*)

hence $p0^{T^*} * x \sqcap y^T \sqcap 1 \leq p0$

by (*metis inf-le1 a4 conv-dist-inf coreflexive-symmetric inf.absorb2 inf.cobounded2 inf.sup-monoid.add-assoc root-var symmetric-one-closed*)

thus $p0^{T^*} * x \sqcap y^T \sqcap 1 \leq p0^{T^*} * x \sqcap p0 \sqcap 1$

by (*meson inf.le-sup-iff order.refl*)

have $p0^{T^*} * x \sqcap p0 \sqcap 1 \leq y$

by (*metis a4 coreflexive-comp-top-inf-one inf.cobounded1 inf-assoc inf-le2*)

thus $p0^{T^*} * x \sqcap p0 \sqcap 1 \leq p0^{T^*} * x \sqcap y^T \sqcap 1$

by (*smt conv-dist-inf coreflexive-symmetric inf.absorb-iff2 inf.cobounded2 inf.sup-monoid.add-assoc*)

qed

thus *?thesis*

by *simp*

qed

also have $\dots = p0 \sqcap 1$

using 1 **by** (*metis inf.sup-monoid.add-commute inf-sup-distrib1 maddux-3-11-pp*)

finally show $p \sqcap 1 = p0 \sqcap 1$

.

show $fc\ p = fc\ p0$
proof (*rule order.antisym*)
have 2: *univalent* $(p0[p0^{T^*} * x \mapsto y])$
by (*simp add: a1 a2 a3 update-univalent mult-assoc*)
have 3: $-(p0^{T^*} * x) \sqcap p0 \leq (p0[p0^{T^*} * x \mapsto y])^* * (p0[p0^{T^*} * x \mapsto y])^{T^*}$
using *fc-increasing inf.order-trans sup.cobounded2* **by** *blast*
have $p0^{T^*} * x \sqcap p0 \leq (p0^{T^*} \sqcap p0 * x^T) * (x \sqcap p0^* * p0)$
by (*metis conv-involutive conv-star-commute dedekind*)
also have $\dots \leq p0^{T^*} * x \sqcap p0 * x^T * p0^* * p0$
by (*metis comp-associative inf.boundedI inf.cobounded2 inf-le1 mult-isotone*)
also have $\dots \leq p0^{T^*} * x \sqcap top * x^T * p0^*$
using *comp-associative comp-inf.mult-right-isotone mult-isotone*
star.right-plus-below-circ **by** *auto*
also have $\dots = p0^{T^*} * x * x^T * p0^*$
by (*metis a2 symmetric-top-closed vector-covector vector-inf-comp*
vector-mult-closed)
also have $\dots \leq (p0^{T^*} * x * y^T) * (y * x^T * p0^*)$
by (*metis a3 order.antisym comp-inf.top-right-mult-increasing*
conv-involutive dedekind-1 inf.sup-left-divisibility inf.sup-monoid.add-commute
mult-right-isotone surjective-conv-total mult-assoc)
also have $\dots = (p0^{T^*} * x \sqcap y^T) * (y \sqcap x^T * p0^*)$
by (*metis a2 a3 vector-covector vector-inf-comp vector-mult-closed*)
also have $\dots = (p0^{T^*} * x \sqcap y^T) * (p0^{T^*} * x \sqcap y^T)^T$
by (*simp add: conv-dist-comp conv-dist-inf conv-star-commute inf-commute*)
also have $\dots \leq (p0[p0^{T^*} * x \mapsto y])^* * (p0[p0^{T^*} * x \mapsto y])^{T^*}$
by (*meson conv-isotone dual-order.trans mult-isotone star.circ-increasing*
sup.cobounded1)
finally have $p0^{T^*} * x \sqcap p0 \leq (p0[p0^{T^*} * x \mapsto y])^* * (p0[p0^{T^*} * x \mapsto y])^{T^*}$
hence $(p0^{T^*} * x \sqcap p0) \sqcup (-(p0^{T^*} * x) \sqcap p0) \leq (p0[p0^{T^*} * x \mapsto y])^* * (p0[p0^{T^*} * x \mapsto y])^{T^*}$
using 3 *le-supI* **by** *blast*
hence $p0 \leq (p0[p0^{T^*} * x \mapsto y])^* * (p0[p0^{T^*} * x \mapsto y])^{T^*}$
using 1 **by** (*metis inf-commute maddux-3-11-pp*)
hence $fc\ p0 \leq (p0[p0^{T^*} * x \mapsto y])^* * (p0[p0^{T^*} * x \mapsto y])^{T^*}$
using 2 *fc-idempotent fc-isotone* **by** *fastforce*
thus $fc\ p0 \leq fc\ p$
by (*simp add: assms(2)*)
have $((p0^{T^*} * x \sqcap y^T) \sqcup (-(p0^{T^*} * x) \sqcap p0))^* = (-(p0^{T^*} * x) \sqcap p0)^* * ((p0^{T^*} * x \sqcap y^T) \sqcup 1)$
proof (*rule star-sup-2*)
have 4: *transitive* $(p0^{T^*} * x)$
using a2 *comp-associative mult-right-isotone rectangle-star-rectangle* **by**
auto
have *transitive* (y^T)
by (*metis a3 conv-dist-comp inf.eq-refl mult-assoc*)
thus *transitive* $(p0^{T^*} * x \sqcap y^T)$
using 4 *transitive-inf-closed* **by** *auto*
have 5: $p0^{T^*} * x * (-(p0^{T^*} * x) \sqcap p0) \leq p0^{T^*} * x$

by (*metis a2 mult-right-isotone top-greatest mult-assoc*)
 have $(-(p0^{T*} * x) \sqcap p0)^T * y \leq p0^T * y$
 by (*simp add: conv-dist-inf mult-left-isotone*)
 also have $\dots \leq y$
 using *a1 a4 root-successor-loop* by *auto*
 finally have $y^T * (-(p0^{T*} * x) \sqcap p0) \leq y^T$
 using *conv-dist-comp conv-isotone* by *fastforce*
 thus $(p0^{T*} * x \sqcap y^T) * (-(p0^{T*} * x) \sqcap p0) \leq p0^{T*} * x \sqcap y^T$
 using *5 comp-left-subdist-inf inf-mono order-trans* by *blast*
 qed
 hence $p^* = (-(p0^{T*} * x) \sqcap p0)^* * ((p0^{T*} * x \sqcap y^T) \sqcup 1)$
 by (*simp add: assms(2)*)
 also have $\dots \leq p0^* * ((p0^{T*} * x \sqcap y^T) \sqcup 1)$
 by (*simp add: mult-left-isotone star-isotone*)
 also have $\dots = p0^* * (p0^{T*} * x * y^T \sqcup 1)$
 by (*simp add: a2 a3 vector-covector vector-mult-closed*)
 also have $\dots = p0^* * (p0^{T*} * (x * x^T) * p0^* * (p0 \sqcap 1) \sqcup 1)$
 by (*metis a4 coreflexive-symmetric inf.cobounded2 root-var comp-associative*
conv-dist-comp conv-involutive conv-star-commute)
 also have $\dots \leq p0^* * (p0^{T*} * 1 * p0^* * (p0 \sqcap 1) \sqcup 1)$
 by (*metis a2 mult-left-isotone mult-right-isotone semiring.add-left-mono*
sup-commute)
 also have $\dots = p0^* * (p0^{T*} * (p0 \sqcap 1) \sqcup p0^* * (p0 \sqcap 1) \sqcup 1)$
 by (*simp add: a1 cancel-separate-eq mult-right-dist-sup*)
 also have $\dots = p0^* * ((p0 \sqcap 1) \sqcup p0^* * (p0 \sqcap 1) \sqcup 1)$
 by (*smt univalent-root-successors a1 conv-dist-comp conv-dist-inf*
coreflexive-idempotent coreflexive-symmetric inf.cobounded2 injective-codomain
loop-root root-transitive-successor-loop symmetric-one-closed)
 also have $\dots = p0^* * (p0^* * (p0 \sqcap 1) \sqcup 1)$
 by (*metis inf.sup-left-divisibility inf-commute sup.left-idem sup-commute*
sup-relative-same-increasing)
 also have $\dots \leq p0^* * p0^*$
 by (*metis inf.cobounded2 inf-commute order.refl order-lesseq-imp*
star.circ-mult-upper-bound star.circ-reflexive star.circ-transitive-equal
sup.boundedI sup-monoid.add-commute)
 also have $\dots = p0^*$
 by (*simp add: star.circ-transitive-equal*)
 finally show $fc\ p \leq fc\ p0$
 by (*metis conv-order conv-star-commute mult-isotone*)
 qed
 qed

lemma *update-acyclic-6:*

assumes *disjoint-set-forest* p

and point x

shows *acyclic* $((p[p^{T*} * x] \rightarrow \text{root } p\ x]) - 1)$

using *assms root-point root-successor-loop update-acyclic-2* by *auto*

theorem *path-compression-assign:*

```

VARs p
[ path-compression-precondition p x y  $\wedge$  p0 = p ]
p[pT* * x] := y
[ path-compression-postcondition p x y p0 ]
apply vcg-tc-simp
apply (unfold path-compression-precondition-def
path-compression-postcondition-def)
apply (intro conjI)
subgoal using update-univalent mult-assoc by auto
subgoal using bijjective-regular mapping-regular regular-closed-star
regular-conv-closed regular-mult-closed update-mapping mult-assoc by auto
subgoal using update-acyclic-6 by blast
subgoal by blast
subgoal by blast
subgoal by blast
subgoal by blast
subgoal by blast
subgoal by blast
subgoal by (smt same-root path-compression-exact
path-compression-precondition-def update-univalent vector-mult-closed)
subgoal using path-compression-exact(1) path-compression-precondition-def
by blast
subgoal using path-compression-exact(2) path-compression-precondition-def
by blast
by blast

```

We next look at implementing these updates using a loop.

```

lemma path-compression-1a:
assumes point x
and disjoint-set-forest p
and  $x \neq \text{root } p$  x
shows  $p^{T^+} * x \leq - x$ 
by (meson assms bijjective-regular mapping-regular regular-closed-star
regular-conv-closed regular-mult-closed vector-mult-closed
point-in-vector-or-complement-2 loop-root-2)

```

```

lemma path-compression-1b:
 $x \leq p^{T^*} * x$ 
using mult-left-isotone star.circ-reflexive by fastforce

```

```

lemma path-compression-1:
path-compression-precondition p x y  $\implies$  path-compression-invariant p x y p x
using path-compression-invariant-def path-compression-precondition-def
loop-root path-compression-1a path-compression-1b by auto

```

```

lemma path-compression-2:
path-compression-invariant p x y p0 w  $\wedge$  y  $\neq$  p[[w]]  $\implies$ 
path-compression-invariant (p[w $\mapsto$ y]) x y p0 (p[[w]])  $\wedge$  ((p[w $\mapsto$ y])T* *
(p[[w]])) $\downarrow$  < (pT* * w) $\downarrow$ 

```

proof –

let $?p = p[w \rightarrow y]$

let $?s = \{ z . \text{regular } z \wedge z \leq p^{T^*} * w \}$

let $?t = \{ z . \text{regular } z \wedge z \leq ?p^{T^*} * (p[[w]]) \}$

assume 1: *path-compression-invariant* $p \ x \ y \ p0 \ w \wedge y \neq p[[w]]$

have $i1$: *disjoint-set-forest* p **and** $i2$: *point* x **and** $i3$: *point* y **and** $i4$: $y = \text{root } p \ x$

using 1 *path-compression-invariant-def path-compression-precondition-def* **by** *meson+*

have $i5$: *point* w **and** $i6$: $y \leq p^{T^*} * w$

and $i7$: $w \neq x \rightarrow p[[x]] = y \wedge y \neq x \wedge p^{T^+} * w \leq -x$ **and** $i8$: $p \sqcap 1 = p0 \sqcap 1$ **and** $i9$: *fc* $p = \text{fc } p0$

and $i10$: *root* $p \ w = y$ **and** $i11$: $p[[w]] = p0[[w]]$ **and** $i12$: $p0[p0^{T^*} * x - p0^{T^*} * w \rightarrow y] = p$

using 1 *path-compression-invariant-def* **by** *blast+*

have $i13$: *disjoint-set-forest* $p0$ **and** $i14$: $y = \text{root } p0 \ x$ **and** $i15$: $w \leq p0^{T^*} * x$

using 1 *path-compression-invariant-def* **by** *auto*

have 2: *point* $(p[[w]])$

using $i1 \ i5$ *read-point* **by** *blast*

show *path-compression-invariant* $?p \ x \ y \ p0 \ (p[[w]]) \wedge \text{card } ?t < \text{card } ?s$

proof (*unfold path-compression-invariant-def, intro conjI*)

have 3: *mapping* $?p$

by (*simp add: i1 i3 i5 bijective-regular update-total update-univalent*)

have 4: $w \neq y$

using 1 $i1 \ i4$ *root-successor-loop* **by** *blast*

hence 5: $w \sqcap y = \text{bot}$

by (*simp add: i3 i5 distinct-points*)

hence $y * w^T \leq -1$

using *pseudo-complement schroeder-4-p* **by** *auto*

hence $y * w^T \leq p^{T^*} - 1$

using $i5 \ i6$ *shunt-bijective* **by** *auto*

also have $\dots \leq p^{T^+}$

by (*simp add: star-plus-without-loops*)

finally have 6: $y \leq p^{T^+} * w$

using $i5$ *shunt-bijective* **by** *auto*

have 7: $w * w^T \leq -p^{T^+}$

proof (*rule ccontr*)

assume $\neg w * w^T \leq -p^{T^+}$

hence $w * w^T \leq --p^{T^+}$

using $i5$ *point-arc arc-in-partition* **by** *blast*

hence $w * w^T \leq p^{T^+} \sqcap 1$

using $i1 \ i5$ *mapping-regular regular-conv-closed regular-closed-star regular-mult-closed* **by** *simp*

also have $\dots = ((p^T \sqcap 1) * p^{T^*} \sqcap 1) \sqcup ((p^T - 1) * p^{T^*} \sqcap 1)$

by (*metis comp-inf.mult-right-dist-sup maddux-3-11-pp mult-right-dist-sup regular-one-closed*)

also have $\dots = ((p^T \sqcap 1) * p^{T^*} \sqcap 1) \sqcup ((p - 1)^+ \sqcap 1)^T$

by (*metis conv-complement conv-dist-inf conv-plus-commute equivalence-one-closed reachable-without-loops*)

also have $\dots \leq ((p^T \sqcap 1) * p^{T*} \sqcap 1) \sqcup (-1 \sqcap 1)^T$
by (*metis (no-types, opaque-lifting) i1 sup-right-isotone inf.sup-left-isotone conv-isotone*)
also have $\dots = (p^T \sqcap 1) * p^{T*} \sqcap 1$
by *simp*
also have $\dots \leq (p^T \sqcap 1) * top \sqcap 1$
by (*metis comp-inf.comp-isotone coreflexive-comp-top-inf equivalence-one-closed inf.cobounded1 inf.cobounded2*)
also have $\dots \leq p^T$
by (*simp add: coreflexive-comp-top-inf-one*)
finally have $w * w^T \leq p^T$
by *simp*
hence $w \leq p[[w]]$
using *i5 shunt-bijective by blast*
hence $w = p[[w]]$
using *2 by (metis i5 epm-3 mult-semi-associative)*
thus *False*
using *2 4 i10 loop-root by auto*
qed
hence *8*: $w \sqcap p^{T+} * w = bot$
using *p-antitone-iff pseudo-complement schroeder-4-p by blast*
show $y \leq ?p^{T*} * (p[[w]])$
proof –
have $(w \sqcap y^T)^T * (-w \sqcap p)^{T*} * p^T * w \leq w^T * (-w \sqcap p)^{T*} * p^T * w$
by (*simp add: conv-isotone mult-left-isotone*)
also have $\dots \leq w^T * p^{T*} * p^T * w$
by (*simp add: conv-isotone mult-left-isotone star-isotone mult-right-isotone*)
also have $\dots = w^T * p^{T+} * w$
by (*simp add: star-plus mult-assoc*)
also have $\dots = bot$
using *8 by (smt i5 covector-inf-comp-3 mult-assoc conv-dist-comp conv-star-commute covector-bot-closed equivalence-top-closed inf.le-iff-sup mult-left-isotone)*
finally have $((w \sqcap y^T)^T \sqcup (-w \sqcap p)^T) * (-w \sqcap p)^{T*} * p^T * w \leq (-w \sqcap p)^T * (-w \sqcap p)^{T*} * p^T * w$
by (*simp add: bot-unique mult-right-dist-sup*)
also have $\dots \leq (-w \sqcap p)^{T*} * p^T * w$
by (*simp add: mult-left-isotone star.left-plus-below-circ*)
finally have $?p^T * (-w \sqcap p)^{T*} * p^T * w \leq (-w \sqcap p)^{T*} * p^T * w$
by (*simp add: conv-dist-sup*)
hence $?p^{T*} * p^T * w \leq (-w \sqcap p)^{T*} * p^T * w$
by (*metis comp-associative star.circ-loop-fixpoint star-left-induct sup-commute sup-least sup-left-divisibility*)
hence $w \sqcap ?p^{T*} * p^T * w \leq w \sqcap (-w \sqcap p)^{T*} * p^T * w$
using *inf.sup-right-isotone by blast*
also have $\dots \leq w \sqcap p^{T*} * p^T * w$
using *conv-isotone mult-left-isotone star-isotone inf.sup-right-isotone by simp*
also have $\dots = bot$

using 8 by (*simp add: star-plus*)
finally have $9: w^T * ?p^{T*} * p^T * w = \text{bot}$
by (*smt i5 covector-inf-comp-3 mult-assoc conv-dist-comp*
covector-bot-closed equivalence-top-closed inf.le-iff-sup mult-left-isotone bot-least
inf.absorb1)
have $p^T * ?p^{T*} * p^T * w = ((w \sqcap p)^T \sqcup (-w \sqcap p)^T) * ?p^{T*} * p^T * w$
by (*metis i5 bijective-regular conv-dist-sup inf.sup-monoid.add-commute*
maddux-3-11-pp)
also have $\dots = (w \sqcap p)^T * ?p^{T*} * p^T * w \sqcup (-w \sqcap p)^T * ?p^{T*} * p^T * w$
by (*simp add: mult-right-dist-sup*)
also have $\dots \leq w^T * ?p^{T*} * p^T * w \sqcup (-w \sqcap p)^T * ?p^{T*} * p^T * w$
using *semiring.add-right-mono comp-isotone conv-isotone* **by** *auto*
also have $\dots = (-w \sqcap p)^T * ?p^{T*} * p^T * w$
using 9 by *simp*
also have $\dots \leq ?p^{T+} * p^T * w$
by (*simp add: conv-isotone mult-left-isotone*)
also have $\dots \leq ?p^{T*} * p^T * w$
by (*simp add: comp-isotone star.left-plus-below-circ*)
finally have $p^{T*} * p^T * w \leq ?p^{T*} * p^T * w$
by (*metis comp-associative star.circ-loop-fixpoint star-left-induct*
sup-commute sup-least sup-left-divisibility)
thus $y \leq ?p^{T*} * (p[[w]])$
using 6 by (*simp add: star-simulation-right-equal mult-assoc*)
qed
have 10: acyclic ($?p - 1$)
using *i1 i10 i3 i5 inf-le1 update-acyclic-3* **by** *blast*
have $?p[[p^{T+} * w]] \leq p^{T+} * w$
proof –
have $(w^T \sqcap y) * p^{T+} * w = y \sqcap w^T * p^{T+} * w$
by (*metis i3 inf-vector-comp vector-inf-comp*)
hence $?p[[p^{T+} * w]] = (y \sqcap w^T * p^{T+} * w) \sqcup (-w^T \sqcap p^T) * p^{T+} * w$
by (*simp add: comp-associative conv-complement conv-dist-inf*
conv-dist-sup mult-right-dist-sup)
also have $\dots \leq y \sqcup (-w^T \sqcap p^T) * p^{T+} * w$
using *sup-left-isotone* **by** *auto*
also have $\dots \leq y \sqcup p^T * p^{T+} * w$
using *mult-left-isotone sup-right-isotone* **by** *auto*
also have $\dots \leq y \sqcup p^{T+} * w$
using *semiring.add-left-mono mult-left-isotone mult-right-isotone*
star.left-plus-below-circ **by** *auto*
also have $\dots = p^{T+} * w$
using 6 by (*simp add: sup-absorb2*)
finally show *?thesis*
by *simp*
qed
hence 11: $?p^{T*} * (p[[w]]) \leq p^{T+} * w$
using *star-left-induct* **by** (*simp add: mult-left-isotone*
star.circ-mult-increasing)
hence 12: $?p^{T+} * (p[[w]]) \leq p^{T+} * w$

using *dual-order.trans mult-left-isotone star.left-plus-below-circ* **by** *blast*
have 13: $?p[[x]] = y \wedge y \neq x \wedge ?p^{T+} * (p[[w]]) \leq -x$
proof (*cases w = x*)
 case *True*
 hence $?p[[x]] = (w^T \sqcap y) * w \sqcup (-w^T \sqcap p^T) * w$
 by (*simp add: conv-complement conv-dist-inf conv-dist-sup*
mult-right-dist-sup)
 also have $\dots = (w^T \sqcap y) * w \sqcup p^T * (-w \sqcap w)$
 by (*metis i5 conv-complement covector-inf-comp-3*
inf.sup-monoid.add-commute vector-complement-closed)
 also have $\dots = (w^T \sqcap y) * w$
 by *simp*
 also have $\dots = y * w$
 by (*simp add: i5 covector-inf-comp-3 inf.sup-monoid.add-commute*)
 also have $\dots = y$
 by (*metis i3 i5 comp-associative*)
 finally show *?thesis*
 using 4 8 12 *True pseudo-complement inf.sup-monoid.add-commute*
order.trans **by** *blast*
 next
 case *False*
 have $?p[[x]] = (w^T \sqcap y) * x \sqcup (-w^T \sqcap p^T) * x$
 by (*simp add: conv-complement conv-dist-inf conv-dist-sup*
mult-right-dist-sup)
 also have $\dots = y * (w \sqcap x) \sqcup p^T * (-w \sqcap x)$
 by (*metis i5 conv-complement covector-inf-comp-3 inf-commute*
vector-complement-closed)
 also have $\dots = p^T * (-w \sqcap x)$
 using *i2 i5 False distinct-points* **by** *auto*
 also have $\dots = y$
 using *i2 i5 i7 False distinct-points inf.absorb2 pseudo-complement* **by** *auto*
 finally show *?thesis*
 using 12 *False i7 dual-order.trans* **by** *blast*
qed
thus $p[[w]] \neq x \longrightarrow ?p[[x]] = y \wedge y \neq x \wedge ?p^{T+} * (p[[w]]) \leq -x$
by *simp*
have 14: $?p^{T*} * x = x \sqcup y$
proof (*rule order.antisym*)
 have $?p^T * (x \sqcup y) = y \sqcup ?p^T * y$
 using 13 **by** (*simp add: mult-left-dist-sup*)
 also have $\dots = y \sqcup (w^T \sqcap y) * y \sqcup (-w^T \sqcap p^T) * y$
 by (*simp add: conv-complement conv-dist-inf conv-dist-sup*
mult-right-dist-sup sup-assoc)
 also have $\dots \leq y \sqcup (w^T \sqcap y) * y \sqcup p^T * y$
 using *mult-left-isotone sup-right-isotone* **by** *auto*
 also have $\dots = y \sqcup (w^T \sqcap y) * y$
 using *i1 i10 root-successor-loop sup-commute* **by** *auto*
 also have $\dots \leq y \sqcup y * y$
 using *mult-left-isotone sup-right-isotone* **by** *auto*

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also have ... =  $y$ 
  by (metis i3 comp-associative sup.idem)
also have ...  $\leq x \sqcup y$ 
  by simp
finally show  $?p^{T^*} * x \leq x \sqcup y$ 
  by (simp add: star-left-induct)
next
  show  $x \sqcup y \leq ?p^{T^*} * x$ 
    using 13 by (metis mult-left-isotone star.circ-increasing
star.circ-loop-fixpoint sup.boundedI sup-ge2)
  qed
have 15:  $y = \text{root } ?p \ x$ 
proof -
  have  $(p \sqcap 1) * y = (p \sqcap 1) * (p \sqcap 1) * p^{T^*} * x$ 
    by (simp add: i4 comp-associative root-var)
  also have ... =  $(p \sqcap 1) * p^{T^*} * x$ 
    using coreflexive-idempotent by auto
  finally have 16:  $(p \sqcap 1) * y = y$ 
    by (simp add: i4 root-var)
  have 17:  $(p \sqcap 1) * x \leq y$ 
    by (metis (no-types, lifting) i4 comp-right-one mult-left-isotone
mult-right-isotone star.circ-reflexive root-var)
  have  $\text{root } ?p \ x = (?p \sqcap 1) * (x \sqcup y)$ 
    using 14 by (metis mult-assoc root-var)
  also have ... =  $(w \sqcap y^T \sqcap 1) * (x \sqcup y) \sqcup (-w \sqcap p \sqcap 1) * (x \sqcup y)$ 
    by (simp add: inf-sup-distrib2 semiring.distrib-right)
  also have ... =  $(w \sqcap 1 \sqcap y^T) * (x \sqcup y) \sqcup (-w \sqcap p \sqcap 1) * (x \sqcup y)$ 
    by (simp add: inf.left-commute inf.sup-monoid.add-commute)
  also have ... =  $(w \sqcap 1) * (y \sqcap (x \sqcup y)) \sqcup (-w \sqcap p \sqcap 1) * (x \sqcup y)$ 
    by (simp add: i3 covector-inf-comp-3)
  also have ... =  $(w \sqcap 1) * y \sqcup (-w \sqcap p \sqcap 1) * (x \sqcup y)$ 
    by (simp add: inf.absorb1)
  also have ... =  $(w \sqcap 1 * y) \sqcup (-w \sqcap (p \sqcap 1) * (x \sqcup y))$ 
    by (simp add: i5 inf-assoc vector-complement-closed vector-inf-comp)
  also have ... =  $(w \sqcap y) \sqcup (-w \sqcap ((p \sqcap 1) * x \sqcup y))$ 
    using 16 by (simp add: mult-left-dist-sup)
  also have ... =  $(w \sqcap y) \sqcup (-w \sqcap y)$ 
    using 17 by (simp add: sup.absorb2)
  also have ... =  $y$ 
    using 5 inf.sup-monoid.add-commute le-iff-inf pseudo-complement
sup-monoid.add-0-left by fastforce
  finally show ?thesis
    by simp
qed
show path-compression-precondition  $?p \ x \ y$ 
  using 3 10 15 i2 i3 path-compression-precondition-def by blast
show vector  $(p[[w]])$ 
  using 2 by simp
show injective  $(p[[w]])$ 

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    using 2 by simp
  show surjective (p[[w]])
    using 2 by simp
  have  $w \sqcap p \sqcap 1 \leq w \sqcap w^T \sqcap p$ 
    by (metis inf.boundedE inf.boundedI inf.cobounded1 inf.cobounded2
one-inf-conv)
  also have  $\dots = w * w^T \sqcap p$ 
    by (simp add: i5 vector-covector)
  also have  $\dots \leq -p^{T+} \sqcap p$ 
    using 7 by (simp add: inf.coboundedI2 inf.sup-monoid.add-commute)
  finally have  $w \sqcap p \sqcap 1 = \text{bot}$ 
    by (metis (no-types, opaque-lifting) conv-dist-inf coreflexive-symmetric
inf.absorb1 inf.boundedE inf.cobounded2 pseudo-complement
star.circ-mult-increasing)
  also have  $w \sqcap y^T \sqcap 1 = \text{bot}$ 
    using 5 antisymmetric-bot-closed asymmetric-bot-closed comp-inf.schroeder-2
inf.absorb1 one-inf-conv by fastforce
  finally have  $w \sqcap p \sqcap 1 = w \sqcap y^T \sqcap 1$ 
    by simp
  thus 18:  $?p \sqcap 1 = p0 \sqcap 1$ 
    by (metis i5 i8 bijective-regular inf.sup-monoid.add-commute inf-sup-distrib2
maddux-3-11-pp)
  show 19:  $fc ?p = fc p0$ 
  proof -
    have  $p[[w]] = p^T * (w \sqcap p^* * y)$ 
      by (metis i3 i5 i6 bijective-reverse conv-star-commute inf.absorb1)
    also have  $\dots = p^T * (w \sqcap p^*) * y$ 
      by (simp add: i5 vector-inf-comp mult-assoc)
    also have  $\dots = p^T * ((w \sqcap 1) \sqcup (w \sqcap p) * (-w \sqcap p)^*) * y$ 
      by (simp add: i5 omit-redundant-points)
    also have  $\dots = p^T * (w \sqcap 1) * y \sqcup p^T * (w \sqcap p) * (-w \sqcap p)^* * y$ 
      by (simp add: comp-associative mult-left-dist-sup mult-right-dist-sup)
    also have  $\dots \leq p^T * y \sqcup p^T * (w \sqcap p) * (-w \sqcap p)^* * y$ 
      by (metis semiring.add-right-mono comp-isotone order.eq-iff
inf.cobounded1 inf.sup-monoid.add-commute mult-1-right)
    also have  $\dots = y \sqcup p^T * (w \sqcap p) * (-w \sqcap p)^* * y$ 
      using i1 i4 root-successor-loop by auto
    also have  $\dots \leq y \sqcup p^T * p * (-w \sqcap p)^* * y$ 
      using comp-isotone sup-right-isotone by auto
    also have  $\dots \leq y \sqcup (-w \sqcap p)^* * y$ 
      by (metis i1 comp-associative eq-refl shunt-mapping sup-right-isotone)
    also have  $\dots = (-w \sqcap p)^* * y$ 
      by (metis star.circ-loop-fixpoint sup.left-idem sup-commute)
    finally have 20:  $p[[w]] \leq (-w \sqcap p)^* * y$ 
      by simp
    have  $p^T * (-w \sqcap p)^* * y = p^T * y \sqcup p^T * (-w \sqcap p) * (-w \sqcap p)^* * y$ 
      by (metis comp-associative mult-left-dist-sup star.circ-loop-fixpoint
sup-commute)
    also have  $\dots = y \sqcup p^T * (-w \sqcap p) * (-w \sqcap p)^* * y$ 

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using *i1 i4 root-successor-loop* **by** *auto*
also have $\dots \leq y \sqcup p^T * p * (-w \sqcap p)^* * y$
using *comp-isotone sup-right-isotone* **by** *auto*
also have $\dots \leq y \sqcup (-w \sqcap p)^* * y$
by (*metis i1 comp-associative eq-refl shunt-mapping sup-right-isotone*)
also have $\dots = (-w \sqcap p)^* * y$
by (*metis star.circ-loop-fixpoint sup.left-idem sup-commute*)
finally have *21*: $p^{T*} * p^T * w \leq (-w \sqcap p)^* * y$
using *20* **by** (*simp add: comp-associative star-left-induct*)
have $w^T \sqcap p^T = p^T * (w^T \sqcap 1)$
by (*metis i5 comp-right-one covector-inf-comp-3*)
inf.sup-monoid.add-commute one-inf-conv
also have $\dots \leq p[[w]]$
by (*metis comp-right-subdist-inf inf.boundedE inf.sup-monoid.add-commute*
one-inf-conv)
also have $\dots \leq p^{T*} * p^T * w$
by (*simp add: mult-left-isotone star.circ-mult-increasing-2*)
also have $\dots \leq (-w \sqcap p)^* * y$
using *21* **by** *simp*
finally have $w \sqcap p \leq y^T * (-w \sqcap p)^{T*}$
by (*metis conv-dist-comp conv-dist-inf conv-involutive conv-isotone*
conv-star-commute)
hence $w \sqcap p \leq (w \sqcap y^T) * (-w \sqcap p)^{T*}$
by (*simp add: i5 vector-inf-comp*)
also have $\dots \leq (w \sqcap y^T) * ?p^{T*}$
by (*simp add: conv-isotone mult-right-isotone star-isotone*)
also have $\dots \leq ?p * ?p^{T*}$
by (*simp add: mult-left-isotone*)
also have $\dots \leq fc ?p$
by (*simp add: mult-left-isotone star.circ-increasing*)
finally have *22*: $w \sqcap p \leq fc ?p$
by *simp*
have $-w \sqcap p \leq ?p$
by *simp*
also have $\dots \leq fc ?p$
by (*simp add: fc-increasing*)
finally have $(w \sqcup -w) \sqcap p \leq fc ?p$
using *22* **by** (*simp add: comp-inf.semiring.distrib-left*
inf.sup-monoid.add-commute)
hence $p \leq fc ?p$
by (*metis i5 bijective-regular inf.sup-monoid.add-commute inf-sup-distrib1*
maddux-3-11-pp)
hence *23*: $fc p \leq fc ?p$
using *3 fc-idempotent fc-isotone* **by** *fastforce*
have $?p \leq (w \sqcap y^T) \sqcup p$
using *sup-right-isotone* **by** *auto*
also have $\dots = w * y^T \sqcup p$
by (*simp add: i3 i5 vector-covector*)
also have $\dots \leq p^* \sqcup p$

by (*smt i5 i6 conv-dist-comp conv-involutive conv-isotone
conv-star-commute le-supI shunt-bijective star.circ-increasing sup-absorb1*)
also have $\dots \leq fc\ p$
using *fc-increasing star.circ-back-loop-prefixpoint* **by** *auto*
finally have $fc\ ?p \leq fc\ p$
using *i1 fc-idempotent fc-isotone* **by** *fastforce*
thus *?thesis*
using *23 i9* **by** *auto*
qed
show $?p[[p[[w]]]] = p0[[p[[w]]]]$
proof –
have $?p[[p[[w]]]] = p[[p[[w]]]]$
using *2 4* **by** (*metis i5 i10 loop-root put-get-different*)
also have $\dots = p[[p0[[w]]]]$
by (*simp add: i11*)
also have $\dots = (p0[p0^{T^*} * x - p0^{T^*} * w \longrightarrow y])[[p0[[w]]]]$
using *i12* **by** *auto*
also have $\dots = p0[[p0[[w]]]]$
proof –
have $p0[[w]] \leq -(p0^{T^*} * x - p0^{T^*} * w)$
by (*meson inf.coboundedI2 mult-left-isotone p-antitone p-antitone-iff
star.circ-increasing*)
thus *?thesis*
by (*meson i2 i5 put-get-different-vector vector-complement-closed
vector-inf-closed vector-mult-closed*)
qed
also have $\dots = p0[[p[[w]]]]$
by (*simp add: i11*)
finally show *?thesis*
qed
have *24*: $root\ ?p\ (p[[w]]) = root\ p0\ (p[[w]])$
using *3 18 19 i13 same-root* **by** *blast*
also have $\dots = root\ p0\ (p0[[w]])$
by (*simp add: i11*)
also have *25*: $\dots = root\ p0\ w$
by (*metis i5 i13 conv-involutive forest-components-increasing
mult-left-isotone shunt-bijective injective-mult-closed read-surjective
same-component-same-root*)
finally show *26*: $root\ ?p\ (p[[w]]) = y$
by (*metis i1 i10 i13 i8 i9 same-root*)
thus $p[[w]] \neq y \longrightarrow ?p^{T^+} * (p[[w]]) \leq -(p[[w]])$
using *2 3 10* **by** (*simp add: path-compression-1a*)
show *univalent p0 total p0 acyclic (p0 - 1)*
by (*simp-all add: i13*)
show $y = root\ p0\ x$
by (*simp add: i14*)
show $p[[w]] \leq p0^{T^*} * x$
by (*metis i11 i15 mult-isotone star.circ-increasing star.circ-transitive-equal*)

mult-assoc)
let $?q = p0[p0^{T^*} * x - p0^{T^*} * (p[[w]]) \dashv\rightarrow y]$
show $?q = ?p$
proof –
have 27: $w \sqcup p0^{T^+} * w = p0^{T^*} * w$
using *comp-associative star.circ-loop-fixpoint sup-commute* **by** *auto*
hence 28: $p0^{T^+} * w = p0^{T^*} * w - w$
using 4 24 25 26 **by** (*metis i11 i13 i5 inf.orderE maddux-3-13*)
path-compression-1a)
hence $p0^{T^*} * (p[[w]]) \leq -w$
by (*metis i11 inf-le2 star-plus mult.assoc*)
hence $w \leq -(p0^{T^*} * (p[[w]]))$
by (*simp add: p-antitone-iff*)
hence $w \leq p0^{T^*} * x - p0^{T^*} * (p[[w]])$
by (*simp add: i15*)
hence 29: $?q \sqcap w = ?p \sqcap w$
by (*metis update-inf update-inf-same*)
have 30: $?q \sqcap p0^{T^+} * w = ?p \sqcap p0^{T^+} * w$
proof –
have $?q \sqcap p0^{T^+} * w = p0 \sqcap p0^{T^+} * w$
by (*metis i11 comp-associative inf.cobounded2 p-antitone-iff*
star.circ-plus-same update-inf-different)
also have $\dots = p \sqcap p0^{T^+} * w$
using 28 **by** (*metis i12 inf.cobounded2 inf.sup-monoid.add-assoc*
p-antitone-iff update-inf-different)
also have $\dots = ?p \sqcap p0^{T^+} * w$
using 28 **by** (*simp add: update-inf-different*)
finally show *?thesis*
.

qed
have 31: $?q \sqcap p0^{T^*} * w = ?p \sqcap p0^{T^*} * w$
using 27 29 30 **by** (*metis inf-sup-distrib1*)
have 32: $?q \sqcap (p0^{T^*} * x - p0^{T^*} * w) = ?p \sqcap (p0^{T^*} * x - p0^{T^*} * w)$
proof –
have $p0^{T^*} * x - p0^{T^*} * w \leq p0^{T^*} * x - p0^{T^*} * (p[[w]])$
using 28 **by** (*metis i11 inf.sup-right-isotone mult.semigroup-axioms*
p-antitone-inf star-plus semigroup.assoc)
hence $?q \sqcap (p0^{T^*} * x - p0^{T^*} * w) = y^T \sqcap p0^{T^*} * x \sqcap -(p0^{T^*} * w)$
by (*metis inf-assoc update-inf*)
also have $\dots = p \sqcap (p0^{T^*} * x - p0^{T^*} * w)$
by (*metis i12 inf-assoc update-inf-same*)
also have $\dots = ?p \sqcap (p0^{T^*} * x - p0^{T^*} * w)$
by (*simp add: inf.coboundedI2 p-antitone path-compression-1b inf-assoc*
update-inf-different)
finally show *?thesis*
.

qed
have $p0^{T^*} * w \sqcup (p0^{T^*} * x - p0^{T^*} * w) = p0^{T^*} * x$
proof –

have 33: *regular* ($p0^{T^*} * w$)
using i13 i5 *bijjective-regular mapping-regular regular-closed-star*
regular-conv-closed regular-mult-closed **by** *auto*
have $p0^{T^*} * w \leq p0^{T^*} * x$
by (*metis i15 comp-associative mult-right-isotone*
star.circ-transitive-equal)
hence $p0^{T^*} * w \sqcup (p0^{T^*} * x - p0^{T^*} * w) = p0^{T^*} * x \sqcap (p0^{T^*} * w \sqcup$
 $-(p0^{T^*} * w))$
by (*simp add: comp-inf.semiring.distrib-left inf.absorb2*)
also have $\dots = p0^{T^*} * x$
using 33 **by** (*metis inf-sup-distrib1 maddux-3-11-pp*)
finally show *?thesis*
.

qed
hence 34: $?q \sqcap p0^{T^*} * x = ?p \sqcap p0^{T^*} * x$
using 31 32 **by** (*metis inf-sup-distrib1*)
have 35: *regular* ($p0^{T^*} * x$)
using i13 i2 *bijjective-regular mapping-regular regular-closed-star*
regular-conv-closed regular-mult-closed **by** *auto*
have $-(p0^{T^*} * x) \leq -w$
by (*simp add: i15 p-antitone*)
hence $?q - p0^{T^*} * x = ?p - p0^{T^*} * x$
by (*metis i12 p-antitone-inf update-inf-different*)
thus *?thesis*
using 34 35 **by** (*metis maddux-3-11-pp*)

qed
show *card ?t < card ?s*
proof –
have $?p^T * p^{T^*} * w = (w^T \sqcap y) * p^{T^*} * w \sqcup (-w^T \sqcap p^T) * p^{T^*} * w$
by (*simp add: conv-complement conv-dist-inf conv-dist-sup*
mult-right-dist-sup)
also have $\dots \leq (w^T \sqcap y) * p^{T^*} * w \sqcup p^T * p^{T^*} * w$
using *mult-left-isotone sup-right-isotone* **by** *auto*
also have $\dots \leq (w^T \sqcap y) * p^{T^*} * w \sqcup p^{T^*} * w$
using *mult-left-isotone star.left-plus-below-circ sup-right-isotone* **by** *blast*
also have $\dots \leq y * p^{T^*} * w \sqcup p^{T^*} * w$
using *semiring.add-right-mono mult-left-isotone* **by** *auto*
also have $\dots \leq y * top \sqcup p^{T^*} * w$
by (*simp add: comp-associative le-supI1 mult-right-isotone*)
also have $\dots = p^{T^*} * w$
by (*simp add: i3 i6 sup-absorb2*)
finally have $?p^{T^*} * p^T * w \leq p^{T^*} * w$
using 11 **by** (*metis dual-order.trans star.circ-loop-fixpoint sup-commute*
sup-ge2 mult-assoc)
hence 36: $?t \subseteq ?s$
using *order-lesseq-imp mult-assoc* **by** *auto*
have 37: $w \in ?s$
by (*simp add: i5 bijjective-regular path-compression-1b*)
have 38: $\neg w \in ?t$

```

proof
  assume  $w \in ?t$ 
  hence 39:  $w \leq (?p^T - 1)^* * (p[[w]])$ 
    using reachable-without-loops by auto
  hence  $p[[w]] \leq (?p - 1)^* * w$ 
    using 2 by (smt i5 bijective-reverse conv-star-commute
reachable-without-loops)
  also have  $\dots \leq p^* * w$ 
  proof -
    have  $p^{T^*} * y = y$ 
      using i1 i4 root-transitive-successor-loop by auto
    hence  $y^T * p^* * w = y^T * w$ 
      by (metis conv-dist-comp conv-involutive conv-star-commute)
    also have  $\dots = \text{bot}$ 
      using 5 by (metis i5 inf.idem inf.sup-monoid.add-commute
mult-left-zero schroeder-1 vector-inf-comp)
    finally have 40:  $y^T * p^* * w = \text{bot}$ 
      by simp
    have  $(?p - 1) * p^* * w = (w \sqcap y^T \sqcap -1) * p^* * w \sqcup (-w \sqcap p \sqcap -1) * p^* * w$ 
      by (simp add: comp-inf.mult-right-dist-sup mult-right-dist-sup)
    also have  $\dots \leq (w \sqcap y^T \sqcap -1) * p^* * w \sqcup p * p^* * w$ 
      by (meson inf-le1 inf-le2 mult-left-isotone order-trans sup-right-isotone)
    also have  $\dots \leq (w \sqcap y^T \sqcap -1) * p^* * w \sqcup p^* * w$ 
      using mult-left-isotone star.left-plus-below-circ sup-right-isotone by blast
    also have  $\dots \leq y^T * p^* * w \sqcup p^* * w$ 
      by (meson inf-le1 inf-le2 mult-left-isotone order-trans sup-left-isotone)
    also have  $\dots = p^* * w$ 
      using 40 by simp
    finally show ?thesis
      by (metis comp-associative le-supI star.circ-loop-fixpoint sup-ge2
star-left-induct)
  qed
  finally have  $w \leq p^{T^*} * p^T * w$ 
    using 11 39 reachable-without-loops star-plus by auto
  thus False
    using 4 i1 i10 i5 loop-root-2 star.circ-plus-same by auto
  qed
  show  $\text{card } ?t < \text{card } ?s$ 
    apply (rule psubset-card-mono)
    subgoal using finite-regular by simp
    subgoal using 36 37 38 by auto
  done
qed
qed
qed

```

lemma *path-compression-3a*:

assumes *path-compression-invariant* p x $(p[[w]])$ $p0$ w

shows $p0[p0^{T^*} * x \rightarrow p[[w]]] = p$
proof –
let $?y = p[[w]]$
let $?p = p0[p0^{T^*} * x \rightarrow ?y]$
have $i1$: *disjoint-set-forest* p **and** $i2$: *point* x **and** $i3$: *point* $?y$ **and** $i4$: $?y = \text{root } p \ x$
using *assms path-compression-invariant-def path-compression-precondition-def*
by *meson+*
have $i5$: *point* w **and** $i6$: $?y \leq p^{T^*} * w$
and $i7$: $w \neq x \rightarrow p[[x]] = ?y \wedge ?y \neq x \wedge p^{T^+} * w \leq -x$ **and** $i8$: $p \sqcap 1 = p0 \sqcap 1$ **and** $i9$: *fc* $p = \text{fc } p0$
and $i10$: *root* $p \ w = ?y$ **and** $i11$: $p[[w]] = p0[[w]]$ **and** $i12$: $p0[p0^{T^*} * x - p0^{T^*} * w \rightarrow ?y] = p$
and $i13$: *disjoint-set-forest* $p0$ **and** $i14$: $?y = \text{root } p0 \ x$ **and** $i15$: $w \leq p0^{T^*} * x$
using *assms path-compression-invariant-def* **by** *blast+*
have 1 : $?p \sqcap ?y = p \sqcap ?y$
by (*metis i1 i14 i3 i4 get-put inf-le1 root-successor-loop update-inf update-inf-same*)
have 2 : $?p \sqcap w = p \sqcap w$
by (*metis i5 i11 i15 get-put update-inf update-inf-same*)
have $?y = \text{root } p0 \ w$
by (*metis i1 i10 i13 i8 i9 same-root*)
hence $p0^{T^*} * w = w \sqcup ?y$
by (*metis i11 i13 root-transitive-successor-loop star.circ-loop-fixpoint star-plus sup-monoid.add-commute mult-assoc*)
hence 3 : $?p \sqcap p0^{T^*} * w = p \sqcap p0^{T^*} * w$
using $1 \ 2$ **by** (*simp add: inf-sup-distrib1*)
have $p0^{T^*} * w \leq p0^{T^*} * x$
by (*metis i15 comp-associative mult-right-isotone star.circ-transitive-equal*)
hence 4 : $?p \sqcap (p0^{T^*} * x \sqcap p0^{T^*} * w) = p \sqcap (p0^{T^*} * x \sqcap p0^{T^*} * w)$
using 3 **by** (*simp add: inf.absorb2*)
have 5 : $?p \sqcap (p0^{T^*} * x - p0^{T^*} * w) = p \sqcap (p0^{T^*} * x - p0^{T^*} * w)$
by (*metis i12 inf-le1 update-inf update-inf-same*)
have *regular* $(p0^{T^*} * w)$
using $i13 \ i5$ *bijective-regular mapping-regular regular-closed-star regular-conv-closed regular-mult-closed* **by** *auto*
hence 6 : $?p \sqcap p0^{T^*} * x = p \sqcap p0^{T^*} * x$
using $4 \ 5$ **by** (*smt inf-sup-distrib1 maddux-3-11-pp*)
have 7 : $?p - p0^{T^*} * x = p - p0^{T^*} * x$
by (*smt i12 inf.sup-monoid.add-commute inf-import-p inf-sup-absorb le-iff-inf p-dist-inf update-inf-different inf.idem p-antitone-inf*)
have *regular* $(p0^{T^*} * x)$
using $i13 \ i2$ *bijective-regular mapping-regular regular-closed-star regular-conv-closed regular-mult-closed* **by** *auto*
thus $?p = p$
using $6 \ 7$ **by** (*smt inf-sup-distrib1 maddux-3-11-pp*)
qed

lemma *path-compression-3*:

$path-compression-invariant\ p\ x\ (p[[w]])\ p0\ w \implies path-compression-postcondition\ p\ x\ (p[[w]])\ p0$
using $path-compression-invariant-def\ path-compression-postcondition-def\ path-compression-precondition-def\ path-compression-3a$ **by** $blast$

theorem $path-compression$:

VARs $p\ t\ w$
 $[path-compression-precondition\ p\ x\ y \wedge p0 = p]$
 $w := x;$
WHILE $y \neq p[[w]]$
INV $\{ path-compression-invariant\ p\ x\ y\ p0\ w \}$
VAR $\{ (p^{T*} * w)\downarrow \}$
DO $t := w;$
 $w := p[[w]];$
 $p[t] := y$
OD
 $[path-compression-postcondition\ p\ x\ y\ p0]$
apply $vcg-tc-simp$
apply $(fact\ path-compression-1)$
apply $(fact\ path-compression-2)$
using $path-compression-3$ **by** $auto$

lemma $path-compression-exists$:

$path-compression-precondition\ p\ x\ y \implies \exists p' . path-compression-postcondition\ p'\ x\ y\ p$
using $tc-extract-function\ path-compression$ **by** $blast$

definition $path-compression\ p\ x\ y \equiv (SOME\ p' . path-compression-postcondition\ p'\ x\ y\ p)$

lemma $path-compression-function$:

assumes $path-compression-precondition\ p\ x\ y$
and $p' = path-compression\ p\ x\ y$
shows $path-compression-postcondition\ p'\ x\ y\ p$
by $(metis\ assms\ path-compression-def\ path-compression-exists\ someI)$

4.4 Find-Set with Path Compression

We sequentially combine find-set and path compression. We consider implementations which use the previously derived functions and implementations which unfold their definitions.

theorem $find-set-path-compression$:

VARs $p\ y$
 $[find-set-precondition\ p\ x \wedge p0 = p]$
 $y := find-set\ p\ x;$
 $p := path-compression\ p\ x\ y$
 $[path-compression-postcondition\ p\ x\ y\ p0]$
apply $vcg-tc-simp$
using $find-set-function\ find-set-postcondition-def\ find-set-precondition-def$

path-compression-function path-compression-precondition-def **by** *fastforce*

theorem *find-set-path-compression-1*:

```

  VARS  $p\ t\ w\ y$ 
  [ find-set-precondition  $p\ x \wedge p0 = p$  ]
   $y := \text{find-set } p\ x;$ 
   $w := x;$ 
  WHILE  $y \neq p[[w]]$ 
    INV { path-compression-invariant  $p\ x\ y\ p0\ w$  }
    VAR {  $(p^{T*} * w)\downarrow$  }
    DO  $t := w;$ 
       $w := p[[w]];$ 
       $p[t] := y$ 
    OD
  [ path-compression-postcondition  $p\ x\ y\ p0$  ]
apply vcg-tc-simp
using find-set-function find-set-postcondition-def find-set-precondition-def
path-compression-1 path-compression-precondition-def apply fastforce
apply (fact path-compression-2)
by (fact path-compression-3)

```

theorem *find-set-path-compression-2*:

```

  VARS  $p\ y$ 
  [ find-set-precondition  $p\ x \wedge p0 = p$  ]
   $y := x;$ 
  WHILE  $y \neq p[[y]]$ 
    INV { find-set-invariant  $p\ x\ y \wedge p0 = p$  }
    VAR {  $(p^{T*} * y)\downarrow$  }
    DO  $y := p[[y]]$ 
    OD;
   $p := \text{path-compression } p\ x\ y$ 
  [ path-compression-postcondition  $p\ x\ y\ p0$  ]
apply vcg-tc-simp
apply (fact find-set-1)
apply (fact find-set-2)
by (smt find-set-3 find-set-invariant-def find-set-postcondition-def
find-set-precondition-def path-compression-function
path-compression-precondition-def)

```

theorem *find-set-path-compression-3*:

```

  VARS  $p\ t\ w\ y$ 
  [ find-set-precondition  $p\ x \wedge p0 = p$  ]
   $y := x;$ 
  WHILE  $y \neq p[[y]]$ 
    INV { find-set-invariant  $p\ x\ y \wedge p0 = p$  }
    VAR {  $(p^{T*} * y)\downarrow$  }
    DO  $y := p[[y]]$ 
    OD;
   $w := x;$ 

```

```

WHILE  $y \neq p[[w]]$ 
  INV { path-compression-invariant  $p\ x\ y\ p0\ w$  }
  VAR {  $(p^{T*} * w)\downarrow$  }
  DO  $t := w;$ 
     $w := p[[w]];$ 
     $p[t] := y$ 
  OD
[ path-compression-postcondition  $p\ x\ y\ p0$  ]
apply vcg-tc-simp
  apply (simp add: find-set-1)
  apply (fact find-set-2)
  using find-set-3 find-set-invariant-def find-set-postcondition-def
find-set-precondition-def path-compression-1 path-compression-precondition-def
apply blast
  apply (fact path-compression-2)
  by (fact path-compression-3)

```

Find-set with path compression returns two results: the representative of the tree and the modified disjoint-set forest.

lemma *find-set-path-compression-exists*:
find-set-precondition $p\ x \implies \exists p'\ y . \text{path-compression-postcondition } p'\ x\ y\ p$
using *tc-extract-function find-set-path-compression* **by** *blast*

definition *find-set-path-compression* $p\ x \equiv (\text{SOME } (p',y) . \text{path-compression-postcondition } p'\ x\ y\ p)$

lemma *find-set-path-compression-function*:
assumes *find-set-precondition* $p\ x$
and $(p',y) = \text{find-set-path-compression } p\ x$
shows *path-compression-postcondition* $p'\ x\ y\ p$
proof –
let $?P = \lambda(p',y) . \text{path-compression-postcondition } p'\ x\ y\ p$
have $?P (\text{SOME } z . ?P\ z)$
apply (*unfold some-eq-ex*)
using *assms(1) find-set-path-compression-exists* **by** *simp*
thus *?thesis*
using *assms(2) find-set-path-compression-def* **by** *auto*
qed

We prove that *find-set-path-compression* returns the same representative as *find-set*.

lemma *find-set-path-compression-find-set*:
assumes *find-set-precondition* $p\ x$
shows *find-set* $p\ x = \text{snd } (\text{find-set-path-compression } p\ x)$
proof –
let $?r = \text{find-set } p\ x$
let $?p = \text{fst } (\text{find-set-path-compression } p\ x)$
let $?y = \text{snd } (\text{find-set-path-compression } p\ x)$
have *1: find-set-postcondition* $p\ x\ ?r$

```

  by (simp add: assms find-set-function)
  have path-compression-postcondition ?p x ?y p
    using assms find-set-path-compression-function prod.collapse by blast
  thus ?r = ?y
    using 1 by (smt assms same-root find-set-precondition-def
find-set-postcondition-def path-compression-precondition-def
path-compression-postcondition-def)
qed

```

A weaker postcondition suffices to prove that the two forests have the same semantics; that is, they describe the same disjoint sets and have the same roots.

lemma *find-set-path-compression-path-compression-semantics:*

```

  assumes find-set-precondition p x
  shows fc (path-compression p x (find-set p x)) = fc (fst
(find-set-path-compression p x))
    and path-compression p x (find-set p x)  $\sqcap$  1 = fst (find-set-path-compression p
x)  $\sqcap$  1
  proof -
    let ?r = find-set p x
    let ?q = path-compression p x ?r
    let ?p = fst (find-set-path-compression p x)
    let ?y = snd (find-set-path-compression p x)
    have 1: path-compression-postcondition (path-compression p x ?r) x ?r p
      using assms find-set-function find-set-postcondition-def
find-set-precondition-def path-compression-function
path-compression-precondition-def by auto
    have 2: path-compression-postcondition ?p x ?y p
      using assms find-set-path-compression-function prod.collapse by blast
    show fc ?q = fc ?p
      using 1 2 by (simp add: path-compression-postcondition-def)
    show ?q  $\sqcap$  1 = ?p  $\sqcap$  1
      using 1 2 by (simp add: path-compression-postcondition-def)
  qed

```

With the current, stronger postcondition of path compression describing the precise effect of how links change, we can prove that the two forests are actually equal.

lemma *find-set-path-compression-find-set-pathcompression:*

```

  assumes find-set-precondition p x
  shows path-compression p x (find-set p x) = fst (find-set-path-compression p x)
  proof -
    let ?r = find-set p x
    let ?q = path-compression p x ?r
    let ?p = fst (find-set-path-compression p x)
    let ?y = snd (find-set-path-compression p x)
    have 1: path-compression-postcondition (path-compression p x ?r) x ?r p
      using assms find-set-function find-set-postcondition-def
find-set-precondition-def path-compression-function

```

```

path-compression-precondition-def by auto
  have 2: path-compression-postcondition ?p x ?y p
    using assms find-set-path-compression-function prod.collapse by blast
  have ?r = ?y
    by (simp add: assms find-set-path-compression-find-set)
  thus ?q = ?p
    using 1 2 by (simp add: path-compression-postcondition-def)
qed

```

4.5 Union-Sets

We only consider a naive union-sets operation (without ranks). The semantics is the equivalence closure obtained after adding the link between the two given nodes, which requires those two elements to be in the same set. The implementation uses temporary variable t to store the two results returned by find-set with path compression. The disjoint-set forest, which keeps being updated, is threaded through the sequence of operations.

definition *union-sets-precondition* $p\ x\ y \equiv \text{disjoint-set-forest } p \wedge \text{point } x \wedge \text{point } y$

definition *union-sets-postcondition* $p\ x\ y\ p0 \equiv \text{union-sets-precondition } p\ x\ y \wedge \text{fc } p = \text{wcc } (p0 \sqcup x * y^T)$

lemma *union-sets-1*:

```

assumes union-sets-precondition p0 x y
  and path-compression-postcondition p1 x r p0
  and path-compression-postcondition p2 y s p1
shows union-sets-postcondition (p2[r $\mapsto$ s]) x y p0
proof (unfold union-sets-postcondition-def union-sets-precondition-def, intro conjI)
  let ?p = p2[r $\mapsto$ s]
  have 1: disjoint-set-forest p1  $\wedge$  point r  $\wedge$  r = root p1 x  $\wedge$  p1  $\sqcap$  1 = p0  $\sqcap$  1  $\wedge$ 
    fc p1 = fc p0
    using assms(2) path-compression-precondition-def
    path-compression-postcondition-def by auto
  have 2: disjoint-set-forest p2  $\wedge$  point s  $\wedge$  s = root p2 y  $\wedge$  p2  $\sqcap$  1 = p1  $\sqcap$  1  $\wedge$ 
    fc p2 = fc p1
    using assms(3) path-compression-precondition-def
    path-compression-postcondition-def by auto
  hence 3: fc p2 = fc p0
    using 1 by simp
  show 4: univalent ?p
    using 1 2 update-univalent by blast
  show total ?p
    using 1 2 bijective-regular update-total by blast
  show acyclic (?p - 1)
proof (cases r = s)
  case True
  thus ?thesis

```



```

    using 2 update-acyclic-5 by fastforce
next
case False
hence bot = r  $\sqcap$  s
  using 1 2 distinct-points by blast
also have ... = r  $\sqcap$  p2T* * s
  using 2 by (smt root-transitive-successor-loop)
finally have s  $\sqcap$  p2* * r = bot
  using schroeder-1 conv-star-commute inf.sup-monoid.add-commute by
fastforce
  thus ?thesis
    using 1 2 update-acyclic-4 by blast
qed
show vector x
  using assms(1) by (simp add: union-sets-precondition-def)
show injective x
  using assms(1) by (simp add: union-sets-precondition-def)
show surjective x
  using assms(1) by (simp add: union-sets-precondition-def)
show vector y
  using assms(1) by (simp add: union-sets-precondition-def)
show injective y
  using assms(1) by (simp add: union-sets-precondition-def)
show surjective y
  using assms(1) by (simp add: union-sets-precondition-def)
show fc ?p = wcc (p0  $\sqcup$  x * yT)
proof (rule order.antisym)
  have r = p1[[r]]
    using 1 by (metis root-successor-loop)
  hence r * rT  $\leq$  p1T
    using 1 eq-refl shunt-bijective by blast
  hence r * rT  $\leq$  p1
    using 1 conv-order coreflexive-symmetric by fastforce
  hence r * rT  $\leq$  p1  $\sqcap$  1
    using 1 inf.boundedI by blast
  also have ... = p2  $\sqcap$  1
    using 2 by simp
  finally have r * rT  $\leq$  p2
    by simp
  hence r  $\leq$  p2 * r
    using 1 shunt-bijective by blast
  hence 5: p2[[r]]  $\leq$  r
    using 2 shunt-mapping by blast
  have r  $\sqcap$  p2  $\leq$  r * (top  $\sqcap$  rT * p2)
    using 1 by (metis dedekind-1)
  also have ... = r * rT * p2
    by (simp add: mult-assoc)
  also have ...  $\leq$  r * rT
    using 5 by (metis comp-associative conv-dist-comp conv-involutive

```

conv-order mult-right-isotone)
also have $\dots \leq 1$
using 1 by *blast*
finally have 6: $r \sqcap p2 \leq 1$
by *simp*
have $p0 \leq wcc\ p0$
by (*simp add: star.circ-sub-dist-1*)
also have $\dots = wcc\ p2$
using 3 by (*simp add: star-decompose-1*)
also have 7: $\dots \leq wcc\ ?p$
proof –
have $wcc\ p2 = wcc\ ((-r \sqcap p2) \sqcup (r \sqcap p2))$
using 1 by (*metis bijective-regular inf.sup-monoid.add-commute*
maddux-3-11-pp)
also have $\dots \leq wcc\ ((-r \sqcap p2) \sqcup 1)$
using 6 *wcc-isotone sup-right-isotone* **by** *simp*
also have $\dots = wcc\ (-r \sqcap p2)$
using *wcc-with-loops* **by** *simp*
also have $\dots \leq wcc\ ?p$
using *wcc-isotone sup-ge2* **by** *blast*
finally show *?thesis*
by *simp*
qed
finally have 8: $p0 \leq wcc\ ?p$
by *force*
have $r \leq p1^{T^*} * x$
using 1 by (*metis inf-le1*)
hence 9: $r * x^T \leq p1^{T^*}$
using *assms(1) shunt-bijective union-sets-precondition-def* **by** *blast*
hence $x * r^T \leq p1^*$
using *conv-dist-comp conv-order conv-star-commute* **by** *force*
also have $\dots \leq wcc\ p1$
by (*simp add: star.circ-sub-dist*)
also have $\dots = wcc\ p2$
using 1 2 by (*simp add: fc-wcc*)
also have $\dots \leq wcc\ ?p$
using 7 by *simp*
finally have 10: $x * r^T \leq wcc\ ?p$
by *simp*
have 11: $r * s^T \leq wcc\ ?p$
using 1 2 *star.circ-sub-dist-1 sup-assoc vector-covector* **by** *auto*
have $s \leq p2^{T^*} * y$
using 2 by (*metis inf-le1*)
hence 12: $s * y^T \leq p2^{T^*}$
using *assms(1) shunt-bijective union-sets-precondition-def* **by** *blast*
also have $\dots \leq wcc\ p2$
using *star-isotone sup-ge2* **by** *blast*
also have $\dots \leq wcc\ ?p$
using 7 by *simp*

finally have 13: $s * y^T \leq wcc \ ?p$
 by *simp*
have $x \leq x * r^T * r \wedge y \leq y * s^T * s$
 using 1 2 *shunt-bijective* by *blast*
hence $x * y^T \leq x * r^T * r * (y * s^T * s)^T$
 using *comp-isotone conv-isotone* by *blast*
also have $\dots = x * r^T * r * s^T * s * y^T$
 by (*simp add: comp-associative conv-dist-comp*)
also have $\dots \leq wcc \ ?p * (r * s^T) * (s * y^T)$
 using 10 by (*metis mult-left-isotone mult-assoc*)
also have $\dots \leq wcc \ ?p * wcc \ ?p * (s * y^T)$
 using 11 by (*metis mult-left-isotone mult-right-isotone*)
also have $\dots \leq wcc \ ?p * wcc \ ?p * wcc \ ?p$
 using 13 by (*metis mult-right-isotone*)
also have $\dots = wcc \ ?p$
 by (*simp add: star.circ-transitive-equal*)
finally have $p0 \sqcup x * y^T \leq wcc \ ?p$
 using 8 by *simp*
hence $wcc (p0 \sqcup x * y^T) \leq wcc \ ?p$
 using *wcc-below-wcc* by *simp*
thus $wcc (p0 \sqcup x * y^T) \leq fc \ ?p$
 using 4 *fc-wcc* by *simp*
have $-r \sqcap p2 \leq wcc \ p2$
 by (*simp add: inf.coboundedI2 star.circ-sub-dist-1*)
also have $\dots = wcc \ p0$
 using 3 by (*simp add: star-decompose-1*)
also have $\dots \leq wcc (p0 \sqcup x * y^T)$
 by (*simp add: wcc-isotone*)
finally have 14: $-r \sqcap p2 \leq wcc (p0 \sqcup x * y^T)$
 by *simp*
have $r * x^T \leq wcc \ p1$
 using 9 *inf.order-trans star.circ-sub-dist sup-commute* by *fastforce*
also have $\dots = wcc \ p0$
 using 1 by (*simp add: star-decompose-1*)
also have $\dots \leq wcc (p0 \sqcup x * y^T)$
 by (*simp add: wcc-isotone*)
finally have 15: $r * x^T \leq wcc (p0 \sqcup x * y^T)$
 by *simp*
have 16: $x * y^T \leq wcc (p0 \sqcup x * y^T)$
 using *le-supE star.circ-sub-dist-1* by *blast*
have $y * s^T \leq p2^*$
 using 12 *conv-dist-comp conv-order conv-star-commute* by *fastforce*
also have $\dots \leq wcc \ p2$
 using *star.circ-sub-dist sup-commute* by *fastforce*
also have $\dots = wcc \ p0$
 using 3 by (*simp add: star-decompose-1*)
also have $\dots \leq wcc (p0 \sqcup x * y^T)$
 by (*simp add: wcc-isotone*)
finally have 17: $y * s^T \leq wcc (p0 \sqcup x * y^T)$

by *simp*
 have $r \leq r * x^T * x \wedge s \leq s * y^T * y$
 using *assms(1) shunt-bijective union-sets-precondition-def* by *blast*
 hence $r * s^T \leq r * x^T * x * (s * y^T * y)^T$
 using *comp-isotone conv-isotone* by *blast*
 also have $\dots = r * x^T * x * y^T * y * s^T$
 by (*simp add: comp-associative conv-dist-comp*)
 also have $\dots \leq wcc (p0 \sqcup x * y^T) * (x * y^T) * (y * s^T)$
 using *15* by (*metis mult-left-isotone mult-assoc*)
 also have $\dots \leq wcc (p0 \sqcup x * y^T) * wcc (p0 \sqcup x * y^T) * (y * s^T)$
 using *16* by (*metis mult-left-isotone mult-right-isotone*)
 also have $\dots \leq wcc (p0 \sqcup x * y^T) * wcc (p0 \sqcup x * y^T) * wcc (p0 \sqcup x * y^T)$
 using *17* by (*metis mult-right-isotone*)
 also have $\dots = wcc (p0 \sqcup x * y^T)$
 by (*simp add: star.circ-transitive-equal*)
 finally have $?p \leq wcc (p0 \sqcup x * y^T)$
 using *1 2 14 vector-covector* by *auto*
 hence $wcc ?p \leq wcc (p0 \sqcup x * y^T)$
 using *wcc-below-wcc* by *blast*
 thus *fc* $?p \leq wcc (p0 \sqcup x * y^T)$
 using *4 fc-wcc* by *simp*

qed

qed

theorem *union-sets*:

VARs p r s t
 [*union-sets-precondition p x y \wedge p0 = p*]
t := find-set-path-compression p x;
p := fst t;
r := snd t;
t := find-set-path-compression p y;
p := fst t;
s := snd t;
p[r] := s
 [*union-sets-postcondition p x y p0*]

proof *vcg-tc-simp*

let *?t1 = find-set-path-compression p0 x*
 let *?p1 = fst ?t1*
 let *?r = snd ?t1*
 let *?t2 = find-set-path-compression ?p1 y*
 let *?p2 = fst ?t2*
 let *?s = snd ?t2*
 let *?p = ?p2[?r \longrightarrow ?s]*
 assume *1: union-sets-precondition p0 x y*
 hence *2: path-compression-postcondition ?p1 x ?r p0*
 by (*simp add: find-set-precondition-def union-sets-precondition-def*
find-set-path-compression-function)
 hence *path-compression-postcondition ?p2 y ?s ?p1*
 using *1* by (*meson find-set-precondition-def union-sets-precondition-def*)

find-set-path-compression-function path-compression-postcondition-def
path-compression-precondition-def prod.collapse)
thus *union-sets-postcondition* (?p2[?r \mapsto ?s]) x y p0
using 1 2 **by** (*simp add: union-sets-1*)
qed

lemma *union-sets-exists:*
union-sets-precondition p x y $\implies \exists p' . \text{union-sets-postcondition } p' x y p$
using *tc-extract-function union-sets* **by** *blast*

definition *union-sets* p x y $\equiv (\text{SOME } p' . \text{union-sets-postcondition } p' x y p)$

lemma *union-sets-function:*
assumes *union-sets-precondition* p x y
and $p' = \text{union-sets } p x y$
shows *union-sets-postcondition* p' x y p
by (*metis assms union-sets-def union-sets-exists someI*)

theorem *union-sets-2:*
VARs p r s
[*union-sets-precondition* p x y $\wedge p0 = p$]
r := *find-set* p x;
p := *path-compression* p x r;
s := *find-set* p y;
p := *path-compression* p y s;
p[r] := s
[*union-sets-postcondition* p x y p0]

proof *vcg-tc-simp*
let ?r = *find-set* p0 x
let ?p1 = *path-compression* p0 x ?r
let ?s = *find-set* ?p1 y
let ?p2 = *path-compression* ?p1 y ?s
assume 1: *union-sets-precondition* p0 x y
hence 2: *path-compression-postcondition* ?p1 x ?r p0
using *find-set-function find-set-postcondition-def find-set-precondition-def*
path-compression-function path-compression-precondition-def
union-sets-precondition-def **by** *auto*
hence *path-compression-postcondition* ?p2 y ?s ?p1
using 1 *find-set-function find-set-postcondition-def find-set-precondition-def*
path-compression-function path-compression-precondition-def
union-sets-precondition-def path-compression-postcondition-def **by** *meson*
thus *union-sets-postcondition* (?p2[?r \mapsto ?s]) x y p0
using 1 2 **by** (*simp add: union-sets-1*)
qed

end

end

theory *More-Disjoint-Set-Forests*

imports *Disjoint-Set-Forests*

begin

5 More on Array Access and Disjoint-Set Forests

This section contains further results about directed acyclic graphs and relational array operations.

context *stone-relation-algebra*

begin

[Theorem 6.4](#)

lemma *update-square*:

assumes *point y*

shows $x[y \mapsto x[x[[y]]]] \leq x * x \sqcup x$

proof –

have $x[y \mapsto x[x[[y]]]] = (y \sqcap y^T * x * x) \sqcup (-y \sqcap x)$

by (*simp add: conv-dist-comp*)

also have $\dots \leq (y \sqcap y^T) * x * x \sqcup x$

by (*smt assms inf.eq-refl inf.sup-monoid.add-commute inf-le1 sup-mono vector-inf-comp*)

also have $\dots \leq x * x \sqcup x$

by (*smt (z3) assms comp-associative conv-dist-comp coreflexive-comp-top-inf inf.cobounded2 sup-left-isotone symmetric-top-closed*)

finally show *?thesis*

qed

[Theorem 2.13](#)

lemma *update-ub*:

$x[y \mapsto z] \leq x \sqcup z^T$

by (*meson dual-order.trans inf.cobounded2 le-supI sup.cobounded1 sup-ge2*)

[Theorem 6.7](#)

lemma *update-square-ub*:

$x[y \mapsto (x * x)^T] \leq x \sqcup x * x$

by (*metis conv-involutive update-ub*)

[Theorem 2.14](#)

lemma *update-same-sub*:

assumes $u \sqcap x = u \sqcap z$

and $y \leq u$

and *regular y*

shows $x[y \mapsto z^T] = x$

by (*smt (z3) assms conv-involutive inf.sup-monoid.add-commute inf.sup-relative-same-increasing maddux-3-11-pp*)

Theorem 2.15

lemma *update-point-get*:

point $y \Longrightarrow x[y \mapsto z[[y]]] = x[y \mapsto z^T]$

by (*metis conv-involutive get-put inf-commute update-inf-same*)

Theorem 2.11

lemma *update-bot*:

$x[\text{bot} \mapsto z] = x$

by *simp*

Theorem 2.12

lemma *update-top*:

$x[\text{top} \mapsto z] = z^T$

by *simp*

Theorem 2.6

lemma *update-same*:

assumes *regular u*

shows $(x[y \mapsto z])[u \mapsto z] = x[y \sqcup u \mapsto z]$

proof –

have $(x[y \mapsto z])[u \mapsto z] = (u \sqcap z^T) \sqcup (-u \sqcap y \sqcap z^T) \sqcup (-u \sqcap -y \sqcap x)$

using *inf.sup-monoid.add-assoc inf-sup-distrib1 sup-assoc* **by** *force*

also have $\dots = (u \sqcap z^T) \sqcup (y \sqcap z^T) \sqcup (-(u \sqcup y) \sqcap x)$

by (*metis assms inf-sup-distrib2 maddux-3-21-pp p-dist-sup*)

also have $\dots = x[y \sqcup u \mapsto z]$

using *comp-inf.mult-right-dist-sup sup-commute* **by** *auto*

finally show *?thesis*

qed

lemma *update-same-3*:

assumes *regular u*

and *regular v*

shows $((x[y \mapsto z])[u \mapsto z])[v \mapsto z] = x[y \sqcup u \sqcup v \mapsto z]$

by (*metis assms update-same*)

Theorem 2.7

lemma *update-split*:

assumes *regular w*

shows $x[y \mapsto z] = (x[y - w \mapsto z])[y \sqcap w \mapsto z]$

by (*smt (z3) assms comp-inf.semiring.distrib-left inf.left-commute inf.sup-monoid.add-commute inf-import-p maddux-3-11-pp maddux-3-12 p-dist-inf sup-assoc*)

Theorem 2.8

lemma *update-injective-swap*:

assumes *injective x*
and *point y*
and *injective z*
and *vector z*
shows *injective* $((x[y \mapsto x[[z]]][z \mapsto x[[y]]])$

proof –

have 1: $(z \sqcap y^T * x) * (z \sqcap y^T * x)^T \leq 1$
using *assms(3) injective-inf-closed by auto*
have $(z \sqcap y^T * x) * (-z \sqcap y \sqcap z^T * x)^T \leq (z \sqcap y^T * x) * (y^T \sqcap x^T * z)$
by *(metis conv-dist-comp conv-involutive conv-order inf.boundedE inf.boundedI inf.cobounded1 inf.cobounded2 mult-right-isotone)*
also have $\dots = (z \sqcap z^T * x) * (y^T \sqcap x^T * y)$
by *(smt (z3) assms(2,4) covector-inf-comp-3 inf.left-commute inf.sup-monoid.add-commute comp-associative conv-dist-comp conv-involutive)*
also have $\dots = (z \sqcap z^T) * x * x^T * (y \sqcap y^T)$
by *(smt (z3) assms(2,4) comp-associative inf.sup-monoid.add-commute vector-covector vector-inf-comp)*
also have $\dots \leq x * x^T$
by *(metis assms(2-4) comp-associative comp-right-one coreflexive-comp-top-inf inf.coboundedI2 mult-right-isotone vector-covector)*
also have $\dots \leq 1$
by *(simp add: assms(1))*
finally have 2: $(z \sqcap y^T * x) * (-z \sqcap y \sqcap z^T * x)^T \leq 1$

have $(z \sqcap y^T * x) * (-z \sqcap -y \sqcap x)^T \leq y^T * x * (-y^T \sqcap x^T)$
by *(smt comp-isotone conv-complement conv-dist-inf inf.cobounded2 inf.sup-monoid.add-assoc)*
also have $\dots = y^T * x * x^T \sqcap -y^T$
by *(simp add: inf.commute assms(2) covector-comp-inf vector-conv-compl)*
also have $\dots \leq y^T \sqcap -y^T$
by *(metis assms(1) comp-associative comp-inf.mult-left-isotone comp-isotone comp-right-one mult-sub-right-one)*
finally have 3: $(z \sqcap y^T * x) * (-z \sqcap -y \sqcap x)^T \leq 1$
using *pseudo-complement by fastforce*
have 4: $(-z \sqcap y \sqcap z^T * x) * (z \sqcap y^T * x)^T \leq 1$
using 2 *conv-dist-comp conv-order by force*
have 5: $(-z \sqcap y \sqcap z^T * x) * (-z \sqcap y \sqcap z^T * x)^T \leq 1$
by *(simp add: assms(2) inf-assoc inf-left-commute injective-inf-closed)*
have $(-z \sqcap y \sqcap z^T * x) * (-z \sqcap -y \sqcap x)^T \leq z^T * x * (-z^T \sqcap x^T)$
using *comp-inf.mult-left-isotone comp-isotone conv-complement conv-dist-inf inf.cobounded1 inf.cobounded2 by auto*
also have $\dots = z^T * x * x^T \sqcap -z^T$
by *(metis assms(4) covector-comp-inf inf.sup-monoid.add-commute vector-conv-compl)*
also have $\dots \leq z^T \sqcap -z^T$
by *(metis assms(1) comp-associative comp-inf.mult-left-isotone comp-isotone comp-right-one mult-sub-right-one)*
finally have 6: $(-z \sqcap y \sqcap z^T * x) * (-z \sqcap -y \sqcap x)^T \leq 1$
using *pseudo-complement by fastforce*

have 7: $(-z \sqcap -y \sqcap x) * (z \sqcap y^T * x)^T \leq 1$
using 3 *conv-dist-comp coreflexive-symmetric by fastforce*
have 8: $(-z \sqcap -y \sqcap x) * (-z \sqcap y \sqcap z^T * x)^T \leq 1$
using 6 *conv-dist-comp coreflexive-symmetric by fastforce*
have 9: $(-z \sqcap -y \sqcap x) * (-z \sqcap -y \sqcap x)^T \leq 1$
using *assms(1) inf.sup-monoid.add-commute injective-inf-closed by auto*
have $(x[y \mapsto x[[z]]][z \mapsto x[[y]]]) = (z \sqcap y^T * x) \sqcup (-z \sqcap y \sqcap z^T * x) \sqcup (-z \sqcap -y \sqcap x)$
by *(simp add: comp-inf.comp-left-dist-sup conv-dist-comp inf-assoc sup-monoid.add-assoc)*
hence $((x[y \mapsto x[[z]]][z \mapsto x[[y]]]) * ((x[y \mapsto x[[z]]][z \mapsto x[[y]]])^T) = ((z \sqcap y^T * x) \sqcup (-z \sqcap y \sqcap z^T * x) \sqcup (-z \sqcap -y \sqcap x)) * ((z \sqcap y^T * x)^T \sqcup (-z \sqcap y \sqcap z^T * x)^T \sqcup (-z \sqcap -y \sqcap x)^T)$
by *(simp add: conv-dist-sup)*
also have $\dots = (z \sqcap y^T * x) * ((z \sqcap y^T * x)^T \sqcup (-z \sqcap y \sqcap z^T * x)^T \sqcup (-z \sqcap -y \sqcap x)^T) \sqcup$
 $(-z \sqcap y \sqcap z^T * x) * ((z \sqcap y^T * x)^T \sqcup (-z \sqcap y \sqcap z^T * x)^T \sqcup (-z \sqcap -y \sqcap x)^T) \sqcup$
 $(-z \sqcap -y \sqcap x) * ((z \sqcap y^T * x)^T \sqcup (-z \sqcap y \sqcap z^T * x)^T \sqcup (-z \sqcap -y \sqcap x)^T)$
using *mult-right-dist-sup by auto*
also have $\dots = (z \sqcap y^T * x) * (z \sqcap y^T * x)^T \sqcup (z \sqcap y^T * x) * (-z \sqcap y \sqcap z^T * x)^T \sqcup$
 $(z \sqcap y^T * x) * (-z \sqcap -y \sqcap x)^T \sqcup$
 $(-z \sqcap y \sqcap z^T * x) * (z \sqcap y^T * x)^T \sqcup (-z \sqcap y \sqcap z^T * x) * (-z \sqcap y \sqcap z^T * x)^T \sqcup$
 $(-z \sqcap y \sqcap z^T * x) * (-z \sqcap -y \sqcap x)^T \sqcup$
 $(-z \sqcap -y \sqcap x) * (z \sqcap y^T * x)^T \sqcup (-z \sqcap -y \sqcap x) * (-z \sqcap y \sqcap z^T * x)^T \sqcup$
 $(-z \sqcap -y \sqcap x) * (-z \sqcap -y \sqcap x)^T$
using *mult-left-dist-sup sup.left-commute sup-commute by auto*
also have $\dots \leq 1$
using 1 2 3 4 5 6 7 8 9 *by simp-all*
finally show *?thesis*

qed

lemma *update-injective-swap-2:*

assumes *injective x*

shows *injective ((x[y ↦ x[[bot]]][bot ↦ x[[y]]])*

by *(simp add: assms inf.sup-monoid.add-commute injective-inf-closed)*

[Theorem 2.9](#)

lemma *update-univalent-swap:*

assumes *univalent x*

and *injective y*

and *vector y*

and *injective z*

and *vector z*

shows *univalent ((x[y ↦ x[[z]]][z ↦ x[[y]]])*

by *(simp add: assms read-injective update-univalent)*

[Theorem 2.10](#)

lemma *update-mapping-swap*:

assumes *mapping x*
and *point y*
and *point z*
shows *mapping* $((x[y \mapsto x[[z]]])[z \mapsto x[[y]])$
by (*simp add: assms bijective-regular read-injective read-surjective update-total update-univalent*)

[Theorem 2.16 *mapping-inf-point-arc*](#) has been moved to theory [Relation-Algebras](#) in entry [Stone-Relation-Algebras](#)

end

context *stone-kleene-relation-algebra*

begin

lemma *omit-redundant-points-2*:

assumes *point p*
shows $p \sqcap x^* = (p \sqcap 1) \sqcup (p \sqcap x \sqcap -p^T) * (x \sqcap -p^T)^*$
proof –
let $?p = p \sqcap 1$
let $?np = -p \sqcap 1$
have $1: p \sqcap x^* \sqcap 1 = p \sqcap 1$
by (*metis inf.le-iff-sup inf.left-commute inf.sup-monoid.add-commute star.circ-reflexive*)
have $2: p \sqcap 1 \sqcap -p^T = \text{bot}$
by (*smt (z3) inf-bot-right inf-commute inf-left-commute one-inf-conv p-inf*)
have $p \sqcap x^* \sqcap -1 = p \sqcap x^* \sqcap -p^T$
by (*metis assms antisymmetric-inf-diversity inf.cobounded1 inf.sup-relative-same-increasing vector-covector*)
also have $\dots = (p \sqcap 1 \sqcap -p^T) \sqcup ((p \sqcap x) * (-p \sqcap x)^* \sqcap -p^T)$
by (*simp add: assms omit-redundant-points comp-inf.semiring.distrib-right*)
also have $\dots = (p \sqcap x) * (-p \sqcap x)^* \sqcap -p^T$
using 2 **by** *simp*
also have $\dots = ?p * x * (-p \sqcap x)^* \sqcap -p^T$
by (*metis assms vector-export-comp-unit*)
also have $\dots = ?p * x * (?np * x)^* \sqcap -p^T$
by (*metis assms vector-complement-closed vector-export-comp-unit*)
also have $\dots = ?p * x * (?np * x)^* * ?np$
by (*metis assms conv-complement covector-comp-inf inf.sup-monoid.add-commute mult-1-right one-inf-conv vector-conv-compl*)
also have $\dots = ?p * x * ?np * (x * ?np)^*$
using *star-slide mult-assoc* **by** *auto*
also have $\dots = (?p * x \sqcap -p^T) * (x * ?np)^*$
by (*metis assms conv-complement covector-comp-inf inf.sup-monoid.add-commute mult-1-right one-inf-conv vector-conv-compl*)
also have $\dots = (?p * x \sqcap -p^T) * (x \sqcap -p^T)^*$
by (*metis assms conv-complement covector-comp-inf inf.sup-monoid.add-commute mult-1-right one-inf-conv vector-conv-compl*)
also have $\dots = (p \sqcap x \sqcap -p^T) * (x \sqcap -p^T)^*$

by (*metis assms vector-export-comp-unit*)
finally show *?thesis*
 using 1 by (*metis maddux-3-11-pp regular-one-closed*)
qed

Theorem 5.3

lemma *omit-redundant-points-3*:

assumes *point p*
shows $p \sqcap x^* = (p \sqcap 1) \sqcup (p \sqcap (x \sqcap -p^T)^+)$
by (*simp add: assms inf-assoc vector-inf-comp omit-redundant-points-2*)

Theorem 6.1

lemma *even-odd-root*:

assumes *acyclic (x - 1)*
and *regular x*
and *univalent x*
shows $(x * x)^{T^*} \sqcap x^T * (x * x)^{T^*} = (1 \sqcap x) * ((x * x)^{T^*} \sqcap x^T * (x * x)^{T^*})$
proof –
have 1: *univalent (x * x)*
by (*simp add: assms(3) univalent-mult-closed*)
have $x \sqcap 1 \leq top * (x \sqcap 1)$
by (*simp add: top-left-mult-increasing*)
hence $x \sqcap -(top * (x \sqcap 1)) \leq x - 1$
using *assms(2) p-shunting-swap pp-dist-comp* **by** *auto*
hence $x^* * (x \sqcap -(top * (x \sqcap 1))) \leq (x - 1)^* * (x - 1)$
using *mult-right-isotone reachable-without-loops* **by** *auto*
also have $\dots \leq -1$
by (*simp add: assms(1) star-plus*)
finally have $(x \sqcap -(top * (x \sqcap 1)))^T \leq -x^*$
using *schroeder-4-p* **by** *force*
hence $x^T \sqcap x^* \leq (top * (x \sqcap 1))^T$
by (*smt (z3) assms(2) conv-complement conv-dist-inf p-shunting-swap regular-closed-inf regular-closed-top regular-mult-closed regular-one-closed*)
also have $\dots = (1 \sqcap x) * top$
by (*metis conv-dist-comp conv-dist-inf inf-commute one-inf-conv symmetric-one-closed symmetric-top-closed*)
finally have 2: $(x^T \sqcap x^*) * top \leq (1 \sqcap x) * top$
by (*metis inf.orderE inf.orderI inf-commute inf-vector-comp*)
have $1 \sqcap x^{T^+} \leq (x^T \sqcap 1 * x^*) * x^{T^*}$
by (*metis conv-involutive conv-star-commute dedekind-2 inf-commute*)
also have $\dots \leq (x^T \sqcap x^*) * top$
by (*simp add: mult-right-isotone*)
also have $\dots \leq (1 \sqcap x) * top$
using 2 **by** *simp*
finally have 3: $1 \sqcap x^{T^+} \leq (1 \sqcap x) * top$
have $x^T \sqcap (x^T * x^T)^+ = 1 * x^T \sqcap (x^T * x^T)^* * x^T * x^T$
using *star-plus mult-assoc* **by** *auto*
also have $\dots = (1 \sqcap (x^T * x^T)^* * x^T) * x^T$

using *assms(3) injective-comp-right-dist-inf* **by force**
also have $\dots \leq (1 \sqcap x^{T^*} * x^T) * x^T$
by (*meson comp-inf.mult-right-isotone comp-isotone inf.eq-refl star.circ-square*)
also have $\dots \leq (1 \sqcap x) * top * x^T$
using *3* **by** (*simp add: mult-left-isotone star-plus*)
also have $\dots \leq (1 \sqcap x) * top$
by (*simp add: comp-associative mult-right-isotone*)
finally have $_4: x^T \sqcap (x^T * x^T)^+ \leq (1 \sqcap x) * top$
 \cdot
have $x^T \sqcap (x^T * x^T)^* = (x^T \sqcap 1) \sqcup (x^T \sqcap (x^T * x^T)^+)$
by (*metis inf-sup-distrib1 star-left-unfold-equal*)
also have $\dots \leq (1 \sqcap x) * top$
using *4* **by** (*metis inf.sup-monoid.add-commute le-supI one-inf-conv top-right-mult-increasing*)
finally have $_4: x^T \sqcap (x^T * x^T)^* \leq (1 \sqcap x) * top$
 \cdot
have $x^T \sqcap (x * x)^* \sqcap -1 \leq x^T \sqcap x^* \sqcap -1$
by (*simp add: inf.coboundedI2 inf.sup-monoid.add-commute star.circ-square*)
also have $\dots = (x - 1)^* \sqcap (x - 1)^T$
using *conv-complement conv-dist-inf inf-assoc inf-left-commute reachable-without-loops symmetric-one-closed* **by auto**
also have $\dots = bot$
using *assms(1) acyclic-star-below-complement-1* **by auto**
finally have *5*: $x^T \sqcap (x * x)^* \sqcap -1 = bot$
by (*simp add: le-bot*)
have $x^T \sqcap (x * x)^* = (x^T \sqcap (x * x)^* \sqcap 1) \sqcup (x^T \sqcap (x * x)^* \sqcap -1)$
by (*metis maddux-3-11-pp regular-one-closed*)
also have $\dots = x^T \sqcap (x * x)^* \sqcap 1$
using *5* **by simp**
also have $\dots = x^T \sqcap 1$
by (*metis calculation comp-inf.semiring.distrib-left inf.sup-monoid.add-commute star.circ-transitive-equal star-involutive star-left-unfold-equal sup-inf-absorb*)
finally have $(x^T \sqcap (x * x)^*) \sqcup (x^T \sqcap (x^T * x^T)^*) \leq (1 \sqcap x) * top$
using *4* *inf.sup-monoid.add-commute one-inf-conv top-right-mult-increasing*
by auto
hence $x^T \sqcap ((x * x)^* \sqcup (x * x)^{T^*}) \leq (1 \sqcap x) * top$
by (*simp add: comp-inf.semiring.distrib-left conv-dist-comp*)
hence *6*: $x^T \sqcap (x * x)^{T^*} * (x * x)^* \leq (1 \sqcap x) * top$
using *1* **by** (*simp add: cancel-separate-eq sup-commute*)
have $(x * x)^{T^*} \sqcap x^T * (x * x)^{T^*} \leq (x^T \sqcap (x * x)^{T^*} * (x * x)^*) * (x * x)^{T^*}$
by (*metis conv-involutive conv-star-commute dedekind-2 inf-commute*)
also have $\dots \leq (1 \sqcap x) * top * (x * x)^{T^*}$
using *6* **by** (*simp add: mult-left-isotone*)
also have $\dots = (1 \sqcap x) * top$
by (*simp add: comp-associative star.circ-left-top*)
finally have $(x * x)^{T^*} \sqcap x^T * (x * x)^{T^*} = (x * x)^{T^*} \sqcap x^T * (x * x)^{T^*} \sqcap (1 \sqcap x) * top$

using *inf.order-iff* **by** *auto*
also have $\dots = (1 \sqcap x) * ((x * x)^{T^*} \sqcap x^T * (x * x)^{T^*})$
by (*metis coreflexive-comp-top-inf inf.cobounded1*
inf.sup-monoid.add-commute)
finally show *?thesis*

\cdot
qed

lemma *update-square-plus*:
 $point\ y \implies x[y \mapsto x[[x[[y]]]]] \leq x^+$
by (*meson update-square comp-isotone dual-order.trans le-supI order-refl*
star.circ-increasing star.circ-mult-increasing)

lemma *update-square-ub-plus*:
 $x[y \mapsto (x * x)^T] \leq x^+$
by (*simp add: comp-isotone inf.coboundedI2 star.circ-increasing*
star.circ-mult-increasing)

Theorem 6.2

lemma *acyclic-square*:
assumes *acyclic* $(x - 1)$
shows $x * x \sqcap 1 = x \sqcap 1$
proof (*rule order.antisym*)
have $1 \sqcap x * x = 1 \sqcap ((x - 1) * x \sqcup (x \sqcap 1) * x)$
by (*metis maddux-3-11-pp regular-one-closed semiring.distrib-right*)
also have $\dots \leq 1 \sqcap ((x - 1) * x \sqcup x)$
by (*metis inf.cobounded2 mult-1-left mult-left-isotone inf.sup-right-isotone*
semiring.add-left-mono)
also have $\dots = 1 \sqcap ((x - 1) * (x - 1) \sqcup (x - 1) * (x \sqcap 1) \sqcup x)$
by (*metis maddux-3-11-pp mult-left-dist-sup regular-one-closed*)
also have $\dots \leq (1 \sqcap (x - 1) * (x - 1)) \sqcup (x - 1) * (x \sqcap 1) \sqcup x$
by (*metis inf-le2 inf-sup-distrib1 semiring.add-left-mono*
sup-monoid.add-assoc)
also have $\dots \leq (1 \sqcap (x - 1)^+) \sqcup (x - 1) * (x \sqcap 1) \sqcup x$
by (*metis comp-isotone inf.eq-refl inf.sup-right-isotone star.circ-increasing*
sup-monoid.add-commute sup-right-isotone)
also have $\dots = (x - 1) * (x \sqcap 1) \sqcup x$
by (*metis assms inf.le-iff-sup inf.sup-monoid.add-commute inf-import-p inf-p*
regular-one-closed sup-inf-absorb sup-monoid.add-commute)
also have $\dots = x$
by (*metis comp-isotone inf.cobounded1 inf-le2 mult-1-right sup.absorb2*)
finally show $x * x \sqcap 1 \leq x \sqcap 1$
by (*simp add: inf.sup-monoid.add-commute*)
show $x \sqcap 1 \leq x * x \sqcap 1$
by (*metis coreflexive-idempotent inf-le1 inf-le2 le-infI mult-isotone*)
qed

lemma *diagonal-update-square-aux*:
assumes *acyclic* $(x - 1)$

and point y
shows $1 \sqcap y \sqcap y^T * x * x = 1 \sqcap y \sqcap x$
proof –
have $1: 1 \sqcap y \sqcap x \leq y^T * x * x$
by (*metis comp-isotone coreflexive-idempotent inf.boundedE inf.cobounded1 inf.cobounded2 one-inf-conv*)
have $1 \sqcap y \sqcap y^T * x * x = 1 \sqcap (y \sqcap y^T) * x * x$
by (*simp add: assms(2) inf.sup-monoid.add-assoc vector-inf-comp*)
also have $\dots = 1 \sqcap (y \sqcap 1) * x * x$
by (*metis assms(2) inf.cobounded1 inf.sup-monoid.add-commute inf.sup-same-context one-inf-conv vector-covector*)
also have $\dots \leq 1 \sqcap x * x$
by (*metis comp-left-subdist-inf inf.sup-right-isotone le-infE mult-left-isotone mult-left-one*)
also have $\dots \leq x$
using *assms(1) acyclic-square inf.sup-monoid.add-commute* **by** *auto*
finally show *?thesis*
using 1 **by** (*metis inf.absorb2 inf.left-commute inf.sup-monoid.add-commute*)
qed

Theorem 6.5

lemma diagonal-update-square:
assumes *acyclic* ($x - 1$)
and point y
shows $(x[y \mapsto x[[x[[y]]]]) \sqcap 1 = x \sqcap 1$
proof –
let $?xy = x[[y]]$
let $?xxy = x[[?xy]]$
let $?xyxxy = x[y \mapsto ?xxy]$
have $?xyxxy \sqcap 1 = ((y \sqcap y^T * x * x) \sqcup (-y \sqcap x)) \sqcap 1$
by (*simp add: conv-dist-comp*)
also have $\dots = (y \sqcap y^T * x * x \sqcap 1) \sqcup (-y \sqcap x \sqcap 1)$
by (*simp add: inf-sup-distrib2*)
also have $\dots = (y \sqcap x \sqcap 1) \sqcup (-y \sqcap x \sqcap 1)$
using *assms* **by** (*smt (verit, ccfv-threshold) diagonal-update-square-aux find-set-precondition-def inf-assoc inf-commute*)
also have $\dots = x \sqcap 1$
by (*metis assms(2) bijective-regular comp-inf.mult-right-dist-sup inf.sup-monoid.add-commute maddux-3-11-pp*)
finally show *?thesis*
qed

Theorem 6.6

lemma fc-update-square:
assumes *mapping* x
and point y
shows $fc(x[y \mapsto x[[x[[y]]]]) = fc\ x$
proof (*rule order.antisym*)

```

let ?xy = x[[y]]
let ?xxy = x[[?xy]]
let ?xyxxy = x[y→?xxy]
have 1: y ⊓ yT * x * x ≤ x * x
  by (smt (z3) assms(2) inf.cobounded2 inf.sup-monoid.add-commute
inf.sup-same-context mult-1-left one-inf-conv vector-covector vector-inf-comp)
have 2: ?xyxxy = (y ⊓ yT * x * x) ⊔ (-y ⊓ x)
  by (simp add: conv-dist-comp)
also have ... ≤ x * x ⊔ x
  using 1 inf-le2 sup-mono by blast
also have ... ≤ x*
  by (simp add: star.circ-increasing star.circ-mult-upper-bound)
finally show fc ?xyxxy ≤ fc x
  by (metis comp-isotone conv-order conv-star-commute star-involutive
star-isotone)
have 3: y ⊓ x ⊓ 1 ≤ fc ?xyxxy
  using inf.coboundedI1 inf.sup-monoid.add-commute reflexive-mult-closed
star.circ-reflexive by auto
have 4: y - 1 ≤ -yT
  using assms(2) p-shunting-swap regular-one-closed vector-covector by auto
have y ⊓ x ≤ yT * x
  by (simp add: assms(2) vector-restrict-comp-conv)
also have ... ≤ yT * x * x * xT
  by (metis assms(1) comp-associative mult-1-right mult-right-isotone total-var)
finally have y ⊓ x ⊓ -1 ≤ y ⊓ -yT ⊓ yT * x * x * xT
  using 4 by (smt (z3) inf.cobounded1 inf.coboundedI2
inf.sup-monoid.add-assoc inf.sup-monoid.add-commute inf-greatest)
also have ... = (y ⊓ yT * x * x) * xT ⊓ -yT
  by (metis assms(2) inf.sup-monoid.add-assoc inf.sup-monoid.add-commute
vector-inf-comp)
also have ... = (y ⊓ yT * x * x) * (xT ⊓ -yT)
  using assms(2) covector-comp-inf vector-conv-compl by auto
also have ... = (y ⊓ yT * x * x) * (-y ⊓ x)T
  by (simp add: conv-complement conv-dist-inf inf-commute)
also have ... ≤ ?xyxxy * (-y ⊓ x)T
  using 2 by (simp add: comp-left-increasing-sup)
also have ... ≤ ?xyxxy * ?xyxxyT
  by (simp add: conv-isotone mult-right-isotone)
also have ... ≤ fc ?xyxxy
  using comp-isotone star.circ-increasing by blast
finally have 5: y ⊓ x ≤ fc ?xyxxy
  using 3 by (smt (z3) comp-inf.semiring.distrib-left inf.le-iff-sup
maddux-3-11-pp regular-one-closed)
have x = (y ⊓ x) ⊔ (-y ⊓ x)
  by (metis assms(2) bijective-regular inf.sup-monoid.add-commute
maddux-3-11-pp)
also have ... ≤ fc ?xyxxy
  using 5 dual-order.trans fc-increasing sup.cobounded2 sup-least by blast
finally show fc x ≤ fc ?xyxxy

```

by (smt (z3) assms fc-equivalence fc-isotone fc-wcc read-injective
star.circ-decompose-9 star-decompose-1 update-univalent)

qed

Theorem 6.2

lemma *acyclic-plus-loop*:

assumes *acyclic* $(x - 1)$

shows $x^+ \sqcap 1 = x \sqcap 1$

proof –

let $?r = x \sqcap 1$

let $?i = x - 1$

have $x^+ \sqcap 1 = (?i \sqcup ?r)^+ \sqcap 1$

by (*metis maddux-3-11-pp regular-one-closed*)

also have $\dots = ((?i^* * ?r)^* * ?i^+ \sqcup (?i^* * ?r)^+) \sqcap 1$

using *plus-sup* by *auto*

also have $\dots \leq (?i^+ \sqcup (?i^* * ?r)^+) \sqcap 1$

by (*metis comp-associative dual-order.eq-iff maddux-3-11-pp
reachable-without-loops regular-one-closed star.circ-plus-same star.circ-sup-9*)

also have $\dots = (?i^* * ?r)^+ \sqcap 1$

by (smt (z3) assms *comp-inf.mult-right-dist-sup inf.absorb2
inf.sup-monoid.add-commute inf-le2 maddux-3-11-pp pseudo-complement
regular-one-closed*)

also have $\dots \leq ?i^* * ?r \sqcap 1$

by (*metis comp-associative dual-order.eq-iff maddux-3-11-pp
reachable-without-loops regular-one-closed star.circ-sup-9 star-slide*)

also have $\dots = (?r \sqcup ?i^+ * ?r) \sqcap 1$

using *comp-associative star.circ-loop-fixpoint sup-commute* by *force*

also have $\dots \leq x \sqcup (?i^+ * ?r \sqcap 1)$

by (*metis comp-inf.mult-right-dist-sup inf.absorb1 inf.cobounded1
inf.cobounded2*)

also have $\dots \leq x \sqcup (-1 * ?r \sqcap 1)$

by (*meson assms comp-inf.comp-isotone mult-left-isotone order.refl
semiring.add-left-mono*)

also have $\dots = x$

by (*metis comp-inf.semiring.mult-not-zero comp-right-one inf.cobounded2
inf-sup-absorb mult-right-isotone pseudo-complement sup.idem sup-inf-distrib1*)

finally show *?thesis*

by (*meson inf.sup-same-context inf-le1 order-trans star.circ-mult-increasing*)

qed

lemma *star-irreflexive-part-eq*:

$x^* - 1 = (x - 1)^+ - 1$

by (*metis reachable-without-loops star-plus-without-loops*)

Theorem 6.3

lemma *star-irreflexive-part*:

$x^* - 1 \leq (x - 1)^+$

using *star-irreflexive-part-eq* by *auto*

lemma *square-irreflexive-part:*

$$x * x - 1 \leq (x - 1)^+$$

proof –

have $x * x = (x \sqcap 1) * x \sqcup (x - 1) * x$

by (*metis maddux-3-11-pp mult-right-dist-sup regular-one-closed*)

also have $\dots \leq 1 * x \sqcup (x - 1) * x$

using *comp-isotone inf.cobounded2 semiring.add-right-mono* **by** *blast*

also have $\dots \leq 1 \sqcup (x - 1) \sqcup (x - 1) * x$

by (*metis inf.cobounded2 maddux-3-11-pp mult-1-left regular-one-closed sup-left-isotone*)

also have $\dots = (x - 1) * (x \sqcup 1) \sqcup 1$

by (*simp add: mult-left-dist-sup sup-assoc sup-commute*)

finally have $x * x - 1 \leq (x - 1) * (x \sqcup 1)$

using *shunting-var-p* **by** *auto*

also have $\dots = (x - 1) * (x - 1) \sqcup (x - 1)$

by (*metis comp-right-one inf.sup-monoid.add-commute maddux-3-21-pp mult-left-dist-sup regular-one-closed sup-commute*)

also have $\dots \leq (x - 1)^+$

by (*metis mult-left-isotone star.circ-increasing star.circ-mult-increasing star.circ-plus-same sup.bounded-iff*)

finally show *?thesis*

qed

Theorem 6.3

lemma *square-irreflexive-part-2:*

$$x * x - 1 \leq x^* - 1$$

using *comp-inf.mult-left-isotone star.circ-increasing star.circ-mult-upper-bound*
by *blast*

Theorem 6.8

lemma *acyclic-update-square:*

assumes *acyclic (x - 1)*

shows *acyclic ((x[y→(x * x)^T]) - 1)*

proof –

have $((x[y \rightarrow (x * x)^T]) - 1)^+ \leq ((x \sqcup x * x) - 1)^+$

by (*metis comp-inf.mult-right-isotone comp-isotone inf.sup-monoid.add-commute star-isotone update-square-ub*)

also have $\dots = ((x - 1) \sqcup (x * x - 1))^+$

using *comp-inf.semiring.distrib-right* **by** *auto*

also have $\dots \leq ((x - 1)^+)^+$

by (*smt (verit, del-insts) comp-isotone reachable-without-loops star.circ-mult-increasing star.circ-plus-same star.circ-right-slide star.circ-separate-5 star.circ-square star.circ-transitive-equal star.left-plus-circ sup.bounded-iff sup-ge1 square-irreflexive-part*)

also have $\dots \leq -1$

using *assms* **by** (*simp add: acyclic-plus*)

finally show *?thesis*

qed

Theorem 6.9

lemma *disjoint-set-forest-update-square*:

assumes *disjoint-set-forest* x

and *vector* y

and *regular* y

shows *disjoint-set-forest* $(x[y \mapsto (x * x)^T])$

proof (*intro conjI*)

show *univalent* $(x[y \mapsto (x * x)^T])$

using *assms update-univalent mapping-mult-closed univalent-conv-injective* **by** *blast*

show *total* $(x[y \mapsto (x * x)^T])$

using *assms update-total total-conv-surjective total-mult-closed* **by** *blast*

show *acyclic* $((x[y \mapsto (x * x)^T]) - 1)$

using *acyclic-update-square assms(1)* **by** *blast*

qed

lemma *disjoint-set-forest-update-square-point*:

assumes *disjoint-set-forest* x

and *point* y

shows *disjoint-set-forest* $(x[y \mapsto (x * x)^T])$

using *assms disjoint-set-forest-update-square bijective-regular* **by** *blast*

end

6 Verifying Further Operations on Disjoint-Set Forests

In this section we verify the *init-sets*, *path-halving* and *path-splitting* operations of disjoint-set forests.

class *choose-point* =

fixes *choose-point* :: $'a \Rightarrow 'a$

Using the *choose-point* operation we define a simple for-each-loop abstraction as syntactic sugar translated to a while-loop. Regular vector h describes the set of all elements that are yet to be processed. It is made explicit so that the invariant can refer to it.

syntax

-Foreach :: $idt \Rightarrow idt \Rightarrow 'assn \Rightarrow 'com \Rightarrow 'com$ $((1FOREACH -/ USING -/ INV \{-\} //DO - /OD) [0,0,0,0] 61)$

translations *FOREACH* x *USING* h *INV* $\{i\}$ *DO* c *OD* \Rightarrow

$h := CONST top;$

WHILE $h \neq CONST bot$

INV $\{CONST regular h \wedge CONST vector h \wedge i\}$

VAR $\{h \downarrow\}$

DO $x := CONST choose-point h;$

```

    c;
    h[x] := CONST bot
  OD

```

```

class stone-kleene-relation-algebra-choose-point-finite-regular =
  stone-kleene-relation-algebra + finite-regular-p-algebra + choose-point +
    assumes choose-point-point: vector  $x \implies x \neq \text{bot} \implies \text{point } (\text{choose-point } x)$ 
    assumes choose-point-decreasing: choose-point  $x \leq --x$ 
begin

```

```

subclass stone-kleene-relation-algebra-tarski-finite-regular
proof unfold-locales

```

```

  fix x
  let ?p = choose-point (x * top)
  let ?q = choose-point ((?p  $\sqcap$  x)T * top)
  let ?y = ?p  $\sqcap$  ?qT
  assume 1: regular x  $x \neq \text{bot}$ 
  hence 2:  $x * \text{top} \neq \text{bot}$ 
    using le-bot top-right-mult-increasing by auto
  hence 3: point ?p
    by (simp add: choose-point-point comp-associative)
  hence 4: ?p  $\neq \text{bot}$ 
    using 2 mult-right-zero by force
  have ?p  $\sqcap$  x  $\neq \text{bot}$ 
proof
  assume ?p  $\sqcap$  x = bot
  hence 5:  $x \leq -?p$ 
    using p-antitone-iff pseudo-complement by auto
  have ?p  $\leq --(x * \text{top})$ 
    by (simp add: choose-point-decreasing)
  also have ...  $\leq --(-?p * \text{top})$ 
    using 5 by (simp add: comp-isotone pp-isotone)
  also have ... = -?p * top
    using regular-mult-closed by auto
  also have ... = -?p
    using 3 vector-complement-closed by auto
  finally have ?p = bot
    using inf-absorb2 by fastforce
  thus False
    using 4 by auto
qed
  hence (?p  $\sqcap$  x)T * top  $\neq \text{bot}$ 
    by (metis comp-inf.semiring.mult-zero-left comp-right-one
  inf.sup-monoid.add-commute inf-top.left-neutral schroeder-1)
  hence point ?q
    using choose-point-point vector-top-closed mult-assoc by auto
  hence 6: arc ?y
    using 3 by (smt bijective-conv-mapping inf.sup-monoid.add-commute
  mapping-inf-point-arc)

```

```

have ?q ≤ --((?p ⊔ x)T * top)
  by (simp add: choose-point-decreasing)
hence ?y ≤ ?p ⊔ --((?p ⊔ x)T * top)T
  by (metis conv-complement conv-isotone inf.sup-right-isotone)
also have ... = ?p ⊔ --(top * (?p ⊔ x))
  by (simp add: conv-dist-comp)
also have ... = ?p ⊔ top * (?p ⊔ x)
  using 1 3 bijective-regular pp-dist-comp by auto
also have ... = ?p ⊔ ?pT * x
  using 3 by (metis comp-inf-vector conv-dist-comp
inf.sup-monoid.add-commute inf-top-right symmetric-top-closed)
also have ... = (?p ⊔ ?pT) * x
  using 3 by (simp add: vector-inf-comp)
also have ... ≤ 1 * x
  using 3 point-antisymmetric mult-left-isotone by blast
finally have ?y ≤ x
  by simp
thus top * x * top = top
  using 6 by (smt (verit, ccfv-SIG) mult-assoc le-iff-sup mult-left-isotone
semiring.distrib-left sup.orderE top.extremum)
qed

```

6.1 Init-Sets

A disjoint-set forest is initialised by applying *make-set* to each node. We prove that the resulting disjoint-set forest is the identity relation.

theorem *init-sets*:

```

VARS h p x
[ True ]
FOREACH x
  USING h
  INV { p - h = 1 - h }
  DO p := make-set p x
  OD
[ p = 1 ∧ disjoint-set-forest p ∧ h = bot ]

```

proof *vcg-tc-simp*

```

fix h p
let ?x = choose-point h
let ?m = make-set p ?x
assume 1: regular h ∧ vector h ∧ p - h = 1 - h ∧ h ≠ bot
show vector (-?x ⊔ h) ∧
  ?m ⊔ (--?x ⊔ -h) = 1 ⊔ (--?x ⊔ -h) ∧
  card { x . regular x ∧ x ≤ -?x ∧ x ≤ h } < h↓

```

proof (*intro conjI*)

```

show vector (-?x ⊔ h)
  using 1 choose-point-point vector-complement-closed vector-inf-closed by
blast

```

blast

```

have 2: point ?x ∧ regular ?x
  using 1 bijective-regular choose-point-point by blast

```

```

have 3:  $-h \leq -?x$ 
  using choose-point-decreasing p-antitone-iff by auto
have 4:  $?x \sqcap ?m = ?x * ?x^T \wedge -?x \sqcap ?m = -?x \sqcap p$ 
  using 1 choose-point-point make-set-function make-set-postcondition-def by
auto
have  $?m \sqcap (--?x \sqcup -h) = (?m \sqcap ?x) \sqcup (?m - h)$ 
  using 2 comp-inf.comp-left-dist-sup by auto
also have  $\dots = ?x * ?x^T \sqcup (?m \sqcap -?x \sqcap -h)$ 
  using 3 4 by (smt (z3) inf-absorb2 inf-assoc inf-commute)
also have  $\dots = ?x * ?x^T \sqcup (1 - h)$ 
  using 1 3 4 inf-absorb2 inf.sup-monoid.add-assoc inf-commute by auto
also have  $\dots = (1 \sqcap ?x) \sqcup (1 - h)$ 
  using 2 by (metis inf.cobounded2 inf.sup-same-context one-inf-conv
vector-covector)
also have  $\dots = 1 \sqcap (--?x \sqcup -h)$ 
  using 2 comp-inf.semiring.distrib-left by auto
finally show  $?m \sqcap (--?x \sqcup -h) = 1 \sqcap (--?x \sqcup -h)$ 
.
have 5:  $\neg ?x \leq -?x$ 
  using 1 2 by (metis comp-commute-below-diversity conv-order
inf.cobounded2 inf-absorb2 pseudo-complement strict-order-var top.extremum)
have 6:  $?x \leq h$ 
  using 1 by (metis choose-point-decreasing)
show card {  $x \cdot \text{regular } x \wedge x \leq -?x \wedge x \leq h$  } < h↓
  apply (rule psubset-card-mono)
  using finite-regular apply simp
  using 2 5 6 by auto
qed
qed
end

```

6.2 Path Halving

Path halving is a variant of the path compression technique. Similarly to path compression, we implement path halving independently of find-set, using a second while-loop which iterates over the same path to the root. We prove that path halving preserves the equivalence-relational semantics of the disjoint-set forest and also preserves the roots of the component trees. Additionally we prove the exact effect of path halving, which is to replace every other parent pointer with a pointer to the respective grandparent.

context *stone-kleene-relation-algebra-tarski-finite-regular*
begin

definition *path-halving-invariant* $p \ x \ y \ p0 \equiv$
find-set-precondition $p \ x \wedge \text{point } y \wedge y \leq p^{T^*} * x \wedge y \leq (p0 * p0)^{T^*} * x \wedge$
 $p0[(p0 * p0)^{T^*} * x - p0^{T^*} * y \longrightarrow (p0 * p0)^T] = p \wedge$
disjoint-set-forest $p0$

definition *path-halving-postcondition* $p\ x\ y\ p0 \equiv$
path-compression-precondition $p\ x\ y \wedge p \sqcap 1 = p0 \sqcap 1 \wedge fc\ p = fc\ p0 \wedge$
 $p0[(p0 * p0)^{T*} * x \mapsto (p0 * p0)^T] = p$

lemma *path-halving-invariant-aux-1*:

assumes *point* x
and *point* y
and *disjoint-set-forest* $p0$
shows $p0 \leq wcc\ (p0[(p0 * p0)^{T*} * x - p0^{T*} * y \mapsto (p0 * p0)^T])$
proof –
let $?p2 = p0 * p0$
let $?p2t = ?p2^T$
let $?p2ts = ?p2t^*$
let $?px = ?p2ts * x$
let $?py = -(p0^{T*} * y)$
let $?pxy = ?px \sqcap ?py$
let $?p = p0[?pxy \mapsto ?p2t]$
have 1: *regular* $?pxy$
using *assms(1,3) bijective-regular find-set-precondition-def mapping-regular pp-dist-comp regular-closed-star regular-conv-closed path-halving-invariant-def* **by** *auto*
have 2: *vector* $x \wedge$ *vector* $?px \wedge$ *vector* $?py$
using *assms(1,2) find-set-precondition-def vector-complement-closed vector-mult-closed path-halving-invariant-def* **by** *auto*
have 3: $?pxy \sqcap p0 \sqcap -?p2 \leq -?px^T$
proof –
have 4: *injective* $x \wedge$ *univalent* $?p2 \wedge$ *regular* $p0$
using *assms(1,3) find-set-precondition-def mapping-regular univalent-mult-closed path-halving-invariant-def* **by** *auto*
have $?p2^* * p0 \sqcap 1 \leq p0^+ \sqcap 1$
using *comp-inf.mult-left-isotone comp-isotone comp-right-one mult-sub-right-one star.circ-square star-slide* **by** *auto*
also **have** $\dots \leq p0$
using *acyclic-plus-loop assms(3) path-halving-invariant-def* **by** *auto*
finally **have** 5: $?p2^* * p0 \sqcap 1 \leq p0$
hence 6: $?p2ts * (1 - p0) \leq -p0$
by (*smt (verit, ccfv-SIG) conv-star-commute dual-order.trans inf.sup-monoid.add-assoc order.refl p-antitone-iff pseudo-complement schroeder-4-p schroeder-6-p*)
have $?p2t^+ * p0 \sqcap 1 = ?p2ts * p0^T * (p0^T * p0) \sqcap 1$
by (*metis conv-dist-comp star-plus mult-assoc*)
also **have** $\dots \leq ?p2ts * p0^T \sqcap 1$
by (*metis assms(3) comp-inf.mult-left-isotone comp-isotone comp-right-one mult-sub-right-one*)
also **have** $\dots \leq p0$
using 5 **by** (*metis conv-dist-comp conv-star-commute inf-commute one-inf-conv star-slide*)
finally **have** $?p2t^+ * p0 \leq -1 \sqcup p0$

by (*metis regular-one-closed shunting-var-p sup-commute*)
 hence $?p2^+ * (1 - p0) \leq -p0$
 by (*smt (z3) conv-dist-comp conv-star-commute half-shunting*
inf.sup-monoid.add-assoc inf.sup-monoid.add-commute pseudo-complement
schroeder-4-p schroeder-6-p star.circ-plus-same)
 have $(1 \sqcap ?px) * top * (1 \sqcap ?px \sqcap -p0) = ?px \sqcap top * (1 \sqcap ?px \sqcap -p0)$
 using 2 by (*metis inf-commute vector-inf-one-comp mult-assoc*)
 also have $\dots = ?px \sqcap ?px^T * (1 - p0)$
 using 2 by (*smt (verit, ccfv-threshold) covector-inf-comp-3*
inf.sup-monoid.add-assoc inf.sup-monoid.add-commute inf-top.left-neutral)
 also have $\dots = ?px \sqcap x^T * ?p2^+ * (1 - p0)$
 by (*simp add: conv-dist-comp conv-star-commute*)
 also have $\dots = (?px \sqcap x^T) * ?p2^+ * (1 - p0)$
 using 2 *vector-inf-comp* by *auto*
 also have $\dots = ?p2ts * (x * x^T) * ?p2^+ * (1 - p0)$
 using 2 *vector-covector mult-assoc* by *auto*
 also have $\dots \leq ?p2ts * ?p2^+ * (1 - p0)$
 using 4 by (*metis inf.order-lesseq-imp mult-left-isotone*
star.circ-mult-upper-bound star.circ-reflexive)
 also have $\dots = (?p2ts \sqcup ?p2^+) * (1 - p0)$
 using 4 by (*simp add: cancel-separate-eq*)
 also have $\dots = (?p2ts \sqcup ?p2^+) * (1 - p0)$
 by (*metis star.circ-plus-one star-plus-loops sup-assoc sup-commute*)
 also have $\dots \leq -p0$
 using 6 7 by (*simp add: mult-right-dist-sup*)
 finally have $(1 \sqcap ?px)^T * p0 * (1 \sqcap ?px \sqcap -p0)^T \leq bot$
 by (*smt (z3) inf.boundedI inf-p top.extremum triple-schroeder-p*)
 hence 8: $(1 \sqcap ?px) * p0 * (1 \sqcap ?px \sqcap -p0) = bot$
 by (*simp add: coreflexive-inf-closed coreflexive-symmetric le-bot*)
 have $?px \sqcap p0 \sqcap ?px^T = (1 \sqcap ?px) * p0 \sqcap ?px^T$
 using 2 *inf-commute vector-inf-one-comp* by *fastforce*
 also have $\dots = (1 \sqcap ?px) * p0 * (1 \sqcap ?px)$
 using 2 by (*metis comp-inf-vector mult-1-right vector-conv-covector*)
 also have $\dots = (1 \sqcap ?px) * p0 * (1 \sqcap ?px \sqcap p0) \sqcup (1 \sqcap ?px) * p0 * (1 \sqcap$
 $?px \sqcap -p0)$
 using 4 by (*metis maddux-3-11-pp mult-left-dist-sup*)
 also have $\dots = (1 \sqcap ?px) * p0 * (1 \sqcap ?px \sqcap p0)$
 using 8 by *simp*
 also have $\dots \leq ?p2$
 by (*metis comp-isotone coreflexive-comp-top-inf inf.cobounded1*
inf.cobounded2)
 finally have $?px \sqcap p0 \sqcap -?p2 \leq -?px^T$
 using 4 *p-shunting-swap regular-mult-closed* by *fastforce*
 thus *?thesis*
 by (*meson comp-inf.mult-left-isotone dual-order.trans inf.cobounded1*)
 qed
 have $p0 \leq ?p2 * p0^T$
 by (*metis assms(3) comp-associative comp-isotone comp-right-one eq-refl*
total-var)

hence $?pxy \sqcap p0 \sqcap -?p2 \leq ?p2 * p0^T$
 by (*metis inf.coboundedI1 inf.sup-monoid.add-commute*)
 hence $?pxy \sqcap p0 \sqcap -?p2 \leq ?pxy \sqcap ?p2 * p0^T \sqcap -?px^T$
 using 3 by (*meson dual-order.trans inf.boundedI inf.cobounded1*)
 also have $\dots = (?pxy \sqcap ?p2) * p0^T \sqcap -?px^T$
 using 2 *vector-inf-comp* by *auto*
 also have $\dots = (?pxy \sqcap ?p2) * (-?px \sqcap p0)^T$
 using 2 by (*simp add: covector-comp-inf inf.sup-monoid.add-commute*
vector-conv-compl conv-complement conv-dist-inf)
 also have $\dots \leq ?p * (-?px \sqcap p0)^T$
 using *comp-left-increasing-sup* by *auto*
 also have $\dots \leq ?p * ?p^T$
 by (*metis comp-inf.mult-right-isotone comp-isotone conv-isotone inf.eq-refl*
inf.sup-monoid.add-commute le-supI1 p-antitone-inf sup-commute)
 also have $\dots \leq wcc ?p$
 using *star.circ-sub-dist-2* by *auto*
 finally have 9: $?pxy \sqcap p0 \sqcap -?p2 \leq wcc ?p$
 .
 have $p0 = (?pxy \sqcap p0) \sqcup (-?pxy \sqcap p0)$
 using 1 by (*metis inf.sup-monoid.add-commute maddux-3-11-pp*)
 also have $\dots \leq (?pxy \sqcap p0) \sqcup ?p$
 using *sup-right-isotone* by *auto*
 also have $\dots = (?pxy \sqcap p0 \sqcap -?p2) \sqcup (?pxy \sqcap p0 \sqcap ?p2) \sqcup ?p$
 by (*smt (z3) assms(3) maddux-3-11-pp mapping-regular pp-dist-comp*
path-halving-invariant-def)
 also have $\dots \leq (?pxy \sqcap p0 \sqcap -?p2) \sqcup (?pxy \sqcap ?p2) \sqcup ?p$
 by (*meson comp-inf.comp-left-subdist-inf inf.boundedE semiring.add-left-mono*
semiring.add-right-mono)
 also have $\dots = (?pxy \sqcap p0 \sqcap -?p2) \sqcup ?p$
 using *sup-assoc* by *auto*
 also have $\dots \leq wcc ?p \sqcup ?p$
 using 9 *sup-left-isotone* by *blast*
 also have $\dots \leq wcc ?p$
 by (*simp add: star.circ-sub-dist-1*)
 finally show *?thesis*
 .

qed

lemma *path-halving-invariant-aux*:

assumes *path-halving-invariant p x y p0*

shows $p[[y]] = p0[[y]]$

and $p[[p[[y]]]] = p0[[p0[[y]]]]$

and $p[[p[[p[[y]]]]]] = p0[[p0[[p0[[y]]]]]]$

and $p \sqcap 1 = p0 \sqcap 1$

and *fc p = fc p0*

proof –

let $?p2 = p0 * p0$

let $?p2t = ?p2^T$

let $?p2ts = ?p2t^*$


```

let ?px = ?p2ts * x
let ?py = -(p0T* * y)
let ?pxy = ?px  $\sqcap$  ?py
let ?p = p0[?pxy  $\rightarrow$  ?p2t]
have ?p[[y]] = p0[[y]]
  apply (rule put-get-different-vector)
  using assms find-set-precondition-def vector-complement-closed
vector-inf-closed vector-mult-closed path-halving-invariant-def apply force
  by (meson inf.cobounded2 order-lesseq-imp p-antitone-iff path-compression-1b)
thus 1: p[[y]] = p0[[y]]
  using assms path-halving-invariant-def by auto
have ?p[[p0[[y]]]] = p0[[p0[[y]]]]
  apply (rule put-get-different-vector)
  using assms find-set-precondition-def vector-complement-closed
vector-inf-closed vector-mult-closed path-halving-invariant-def apply force
  by (metis comp-isotone inf.boundedE inf.coboundedI2 inf.eq-refl p-antitone-iff
selection-closed-id star.circ-increasing)
thus 2: p[[p[[y]]]] = p0[[p0[[y]]]]
  using 1 assms path-halving-invariant-def by auto
have ?p[[p0[[p0[[y]]]]]] = p0[[p0[[p0[[y]]]]]]
  apply (rule put-get-different-vector)
  using assms find-set-precondition-def vector-complement-closed
vector-inf-closed vector-mult-closed path-halving-invariant-def apply force
  by (metis comp-associative comp-isotone conv-dist-comp conv-involutive
conv-order inf.coboundedI2 inf.le-iff-sup mult-left-isotone p-antitone-iff
p-antitone-inf star.circ-increasing star.circ-transitive-equal)
thus p[[p[[p[[y]]]]]] = p0[[p0[[p0[[y]]]]]]
  using 2 assms path-halving-invariant-def by auto
have 3: regular ?pxy
  using assms bijective-regular find-set-precondition-def mapping-regular
pp-dist-comp regular-closed-star regular-conv-closed path-halving-invariant-def by
auto
have p  $\sqcap$  1 = ?p  $\sqcap$  1
  using assms path-halving-invariant-def by auto
also have ... = (?pxy  $\sqcap$  ?p2  $\sqcap$  1)  $\sqcup$  (-?pxy  $\sqcap$  p0  $\sqcap$  1)
  using comp-inf.semiring.distrib-right conv-involutive by auto
also have ... = (?pxy  $\sqcap$  p0  $\sqcap$  1)  $\sqcup$  (-?pxy  $\sqcap$  p0  $\sqcap$  1)
  using assms acyclic-square path-halving-invariant-def
inf.sup-monoid.add-assoc by auto
also have ... = (?pxy  $\sqcup$  -?pxy)  $\sqcap$  p0  $\sqcap$  1
  using inf-sup-distrib2 by auto
also have ... = p0  $\sqcap$  1
  using 3 by (metis inf.sup-monoid.add-commute inf-sup-distrib1
maddux-3-11-pp)
finally show p  $\sqcap$  1 = p0  $\sqcap$  1
  .
have p  $\leq$  p0+
  by (metis assms path-halving-invariant-def update-square-ub-plus)
hence 4: fc p  $\leq$  fc p0

```

using *conv-plus-commute fc-isotone star.left-plus-circ* **by** *fastforce*
have $wcc\ p0 \leq wcc\ ?p$
by (*meson assms wcc-below-wcc path-halving-invariant-aux-1*
path-halving-invariant-def find-set-precondition-def)
hence $fc\ p0 \leq fc\ ?p$
using *assms find-set-precondition-def path-halving-invariant-def fc-wcc* **by** *auto*
thus $fc\ p = fc\ p0$
using 4 *assms path-halving-invariant-def* **by** *auto*
qed

lemma *path-halving-1*:

find-set-precondition p0 x \implies path-halving-invariant p0 x x p0

proof –

assume 1: *find-set-precondition p0 x*
show *path-halving-invariant p0 x x p0*
proof (*unfold path-halving-invariant-def, intro conjI*)
show *find-set-precondition p0 x*
using 1 **by** *simp*
show *vector x injective x surjective x*
using 1 *find-set-precondition-def* **by** *auto*
show $x \leq p0^{T^*} * x$
by (*simp add: path-compression-1b*)
show $x \leq (p0 * p0)^{T^*} * x$
by (*simp add: path-compression-1b*)
have $(p0 * p0)^{T^*} * x \leq p0^{T^*} * x$
by (*simp add: conv-dist-comp mult-left-isotone star.circ-square*)
thus $p0[(p0 * p0)^{T^*} * x - p0^{T^*} * x \longrightarrow (p0 * p0)^T] = p0$
by (*smt (z3) inf.le-iff-sup inf-commute maddux-3-11-pp p-antitone-inf*
pseudo-complement)
show *univalent p0 total p0 acyclic (p0 - 1)*
using 1 *find-set-precondition-def* **by** *auto*
qed
qed

lemma *path-halving-2*:

path-halving-invariant p x y p0 \wedge y \neq p[[y]] \implies path-halving-invariant
(p[y \mapsto p[[p[[[y]]]]]]) x ((p[y \mapsto p[[p[[[y]]]]]])[[y]]) p0 \wedge ((p[y \mapsto p[[p[[[y]]]]]])^{T} **
((p[y \mapsto p[[p[[[y]]]]]])[[y]]) \downarrow < (p^{T} * y) \downarrow*

proof –

let $?py = p[[y]]$
let $?ppy = p[[?py]]$
let $?pyppy = p[y \mapsto ?ppy]$
let $?p2 = p0 * p0$
let $?p2t = ?p2^T$
let $?p2ts = ?p2t^*$
let $?px = ?p2ts * x$
let $?py2 = -(p0^{T^*} * y)$
let $?pxy = ?px \sqcap ?py2$
let $?p = p0[?pxy \mapsto ?p2t]$

```

let ?pty = p0T * y
let ?pt2y = p0T * p0T * y
let ?pt2sy = p0T* * p0T * p0T * y
assume 1: path-halving-invariant p x y p0 ∧ y ≠ ?py
have 2: point ?pty ∧ point ?pt2y
  using 1 by (smt (verit) comp-associative read-injective read-surjective
path-halving-invariant-def)
show path-halving-invariant ?pyppy x (?pyppy[[y]]) p0 ∧ (?pyppyT* *
(?pyppy[[y]]))↓ < (pT* * y)↓
proof
  show path-halving-invariant ?pyppy x (?pyppy[[y]]) p0
  proof (unfold path-halving-invariant-def, intro conjI)
    show 3: find-set-precondition ?pyppy x
    proof (unfold find-set-precondition-def, intro conjI)
      show univalent ?pyppy
        using 1 find-set-precondition-def read-injective update-univalent
path-halving-invariant-def by auto
      show total ?pyppy
        using 1 bijective-regular find-set-precondition-def read-surjective
update-total path-halving-invariant-def by force
      show acyclic (?pyppy - 1)
        apply (rule update-acyclic-3)
        using 1 find-set-precondition-def path-halving-invariant-def apply blast
        using 1 2 comp-associative path-halving-invariant-aux(2) apply force
        using 1 path-halving-invariant-def apply blast
        by (metis inf.order-lesseq-imp mult-isotone star.circ-increasing
star.circ-square mult-assoc)
      show vector x injective x surjective x
        using 1 find-set-precondition-def path-halving-invariant-def by auto
    qed
    show vector (?pyppy[[y]])
      using 1 comp-associative path-halving-invariant-def by auto
    show injective (?pyppy[[y]])
      using 1 3 read-injective path-halving-invariant-def find-set-precondition-def
by auto
    show surjective (?pyppy[[y]])
      using 1 3 read-surjective path-halving-invariant-def
find-set-precondition-def by auto
    show ?pyppy[[y]] ≤ ?pyppyT* * x
    proof -
      have y = (y ⊓ pT*) * x
        using 1 le-iff-inf vector-inf-comp path-halving-invariant-def by auto
      also have ... = ((y ⊓ 1) ⊔ (y ⊓ (pT ⊓ -yT)+)) * x
        using 1 omit-redundant-points-3 path-halving-invariant-def by auto
      also have ... ≤ (1 ⊔ (y ⊓ (pT ⊓ -yT)+)) * x
        using 1 sup-inf-distrib2 vector-inf-comp path-halving-invariant-def by
auto
      also have ... ≤ (1 ⊔ (pT ⊓ -yT)+) * x
        by (simp add: inf.coboundedI2 mult-left-isotone)

```

```

also have ... = (p  $\sqcap$  -y)T* * x
  by (simp add: conv-complement conv-dist-inf star-left-unfold-equal)
also have ... ≤ ?pyppyT* * x
  by (simp add: conv-isotone inf.sup-monoid.add-commute mult-left-isotone
star-isotone)
  finally show ?thesis
    by (metis mult-isotone star.circ-increasing star.circ-transitive-equal
mult-assoc)
qed
show ?pyppy[[y]] ≤ ?px
proof -
  have ?pyppy[[y]] = p[[?py]]
    using 1 put-get vector-mult-closed path-halving-invariant-def by force
  also have ... = p0[[p0[[y]]]]
    using 1 path-halving-invariant-aux(2) by blast
  also have ... = ?p2t * y
    by (simp add: conv-dist-comp mult-assoc)
  also have ... ≤ ?p2t * ?px
    using 1 path-halving-invariant-def comp-associative mult-right-isotone by
force
  also have ... ≤ ?px
    by (metis comp-associative mult-left-isotone star.left-plus-below-circ)
  finally show ?thesis
    .
qed
show p0[?px - p0T* * (?pyppy[[y]])] → ?p2t = ?pyppy
proof -
  have ?px  $\sqcap$  ?pty = ?px  $\sqcap$  p0T * ?px  $\sqcap$  ?pty
    using 1 inf.absorb2 inf.sup-monoid.add-assoc mult-right-isotone
path-halving-invariant-def by force
  also have ... = (?p2ts  $\sqcap$  p0T * ?p2ts) * x  $\sqcap$  ?pty
    using 3 comp-associative find-set-precondition-def
injective-comp-right-dist-inf by auto
  also have ... = (1  $\sqcap$  p0) * (?p2ts  $\sqcap$  p0T * ?p2ts) * x  $\sqcap$  ?pty
    using 1 even-odd-root mapping-regular path-halving-invariant-def by auto
  also have ... ≤ (1  $\sqcap$  p0) * top  $\sqcap$  ?pty
    by (metis comp-associative comp-inf.mult-left-isotone
comp-inf.star.circ-sub-dist-2 comp-left-subdist-inf dual-order.trans
mult-right-isotone)
  also have 4: ... = (1  $\sqcap$  p0T) * ?pty
    using coreflexive-comp-top-inf one-inf-conv by auto
  also have ... ≤ ?pt2y
    by (simp add: mult-assoc mult-left-isotone)
  finally have 5: ?px  $\sqcap$  ?pty ≤ ?pt2y
    .
  have 6: p[?px  $\sqcap$  -?pt2sy  $\sqcap$  ?pty] → ?p2t = p
proof (cases ?pty ≤ ?px  $\sqcap$  -?pt2sy)
  case True
    hence ?pty ≤ ?pt2y

```

```

    using 5 conv-dist-comp inf.absorb2 by auto
  hence 7: ?pty = ?pt2y
    using 2 epm-3 by fastforce
  have p[?px  $\sqcap$  -?pt2sy  $\sqcap$  ?pty $\rightarrow$ ?p2t] = p[?pty $\rightarrow$ ?p2t]
    using True inf.absorb2 by auto
  also have ... = p[?pty $\rightarrow$ ?p2[[?pty]]]
    using 2 update-point-get by auto
  also have ... = p[?pty $\rightarrow$ p0T * p0T * p0T * y]
    using comp-associative conv-dist-comp by auto
  also have ... = p[?pty $\rightarrow$ ?pt2y]
    using 7 mult-assoc by simp
  also have ... = p[?pty $\rightarrow$ p[[?pty]]]
    using 1 path-halving-invariant-aux(1,2) mult-assoc by force
  also have ... = p
    using 2 get-put by auto
  finally show ?thesis
.
next
case False
have mapping ?p2
  using 1 mapping-mult-closed path-halving-invariant-def by blast
hence 8: regular (?px  $\sqcap$  -?pt2sy)
  using 1 bijective-regular find-set-precondition-def mapping-regular
pp-dist-comp regular-closed-star regular-conv-closed path-halving-invariant-def by
auto
  have vector (?px  $\sqcap$  -?pt2sy)
    using 1 find-set-precondition-def vector-complement-closed
vector-inf-closed vector-mult-closed path-halving-invariant-def by force
  hence ?pty  $\leq$  -(?px  $\sqcap$  -?pt2sy)
    using 2 8 point-in-vector-or-complement False by blast
  hence ?px  $\sqcap$  -?pt2sy  $\sqcap$  ?pty = bot
    by (simp add: p-antitone-iff pseudo-complement)
  thus ?thesis
    by simp
qed
have 9: p[?px  $\sqcap$  -?pt2sy  $\sqcap$  y $\rightarrow$ ?p2t] = ?pyppy
proof (cases y  $\leq$  -?pt2sy)
case True
  hence p[?px  $\sqcap$  -?pt2sy  $\sqcap$  y $\rightarrow$ ?p2t] = p[y $\rightarrow$ ?p2t]
    using 1 inf.absorb2 path-halving-invariant-def by auto
  also have ... = ?pyppy
    using 1 by (metis comp-associative conv-dist-comp
path-halving-invariant-aux(2) path-halving-invariant-def update-point-get)
  finally show ?thesis
.
next
case False
have vector (-?pt2sy)
  using 1 vector-complement-closed vector-mult-closed

```

path-halving-invariant-def **by** *blast*
hence 10: $y \leq ?pt2sy$
using 1 **by** (*smt* (*verit*, *del-insts*) *False* *bijjective-regular*
point-in-vector-or-complement *regular-closed-star* *regular-mult-closed*
total-conv-surjective *univalent-conv-injective* *path-halving-invariant-def*)
hence $?px \sqcap -?pt2sy \sqcap y = \text{bot}$
by (*simp* *add*: *inf.coboundedI2* *p-antitone* *pseudo-complement*)
hence 11: $p[?px \sqcap -?pt2sy \sqcap y \multimap ?p2t] = p$
by *simp*
have $y \leq p0^{T+} * y$
using 10 **by** (*metis* *mult-left-isotone* *order-lesseq-imp*
star.circ-plus-same *star.left-plus-below-circ*)
hence 12: $y = \text{root } p0 \ y$
using 1 *loop-root-2* *path-halving-invariant-def* **by** *blast*
have $?pyppy = p[y \multimap p0[[p0[[y]]]]]$
using 1 *path-halving-invariant-aux*(2) **by** *force*
also have $\dots = p[y \multimap p0[[y]]]$
using 1 12 **by** (*metis* *root-successor-loop* *path-halving-invariant-def*)
also have $\dots = p[y \multimap ?py]$
using 1 *path-halving-invariant-aux*(1) **by** *force*
also have $\dots = p$
using 1 *get-put* *path-halving-invariant-def* **by** *blast*
finally show *?thesis*
using 11 **by** *simp*
qed
have 13: $-?pt2sy = -(p0^{T*} * y) \sqcup (-?pt2sy \sqcap ?pty) \sqcup (-?pt2sy \sqcap y)$
proof (*rule* *order.antisym*)
have 14: *regular* ($p0^{T*} * y$) \wedge *regular* $?pt2sy$
using 1 **by** (*metis* *order.antisym* *conv-complement* *conv-dist-comp*
conv-involutive *conv-star-commute* *forest-components-increasing* *mapping-regular*
pp-dist-star *regular-mult-closed* *top.extremum* *path-halving-invariant-def*)
have $p0^{T*} = p0^{T*} * p0^T \sqcup p0^T \sqcup 1$
using *star.circ-back-loop-fixpoint* *star.circ-plus-same*
star-left-unfold-equal *sup-commute* **by** *auto*
hence $p0^{T*} * y \leq ?pt2sy \sqcup ?pty \sqcup y$
by (*metis* *inf.eq-refl* *mult-1-left* *mult-right-dist-sup*)
also have $\dots = ?pt2sy \sqcup (-?pt2sy \sqcap ?pty) \sqcup y$
using 14 **by** (*metis* *maddux-3-21-pp*)
also have $\dots = ?pt2sy \sqcup (-?pt2sy \sqcap ?pty) \sqcup (-?pt2sy \sqcap y)$
using 14 **by** (*smt* (*z3*) *maddux-3-21-pp* *sup.left-commute* *sup-assoc*)
hence $p0^{T*} * y \sqcap -?pt2sy \leq (-?pt2sy \sqcap ?pty) \sqcup (-?pt2sy \sqcap y)$
using *calculation* *half-shunting* *sup-assoc* *sup-commute* **by** *auto*
thus $-?pt2sy \leq -(p0^{T*} * y) \sqcup (-?pt2sy \sqcap ?pty) \sqcup (-?pt2sy \sqcap y)$
using 14 **by** (*smt* (*z3*) *inf.sup-monoid.add-commute* *shunting-var-p*
sup.left-commute *sup-commute*)
have $-(p0^{T*} * y) \leq -?pt2sy$
by (*meson* *mult-left-isotone* *order.trans* *p-antitone*
star.right-plus-below-circ)
thus $-(p0^{T*} * y) \sqcup (-?pt2sy \sqcap ?pty) \sqcup (-?pt2sy \sqcap y) \leq -?pt2sy$

by *simp*
 qed
 have *regular* ?*px* *regular* ?*pty* *regular* *y*
 using 1 *bijjective-regular find-set-precondition-def mapping-regular*
pp-dist-comp regular-closed-star regular-conv-closed path-halving-invariant-def by *auto*
 hence 15: *regular* (?*px* \sqcap - ?*pt2sy* \sqcap ?*pty*) *regular* (?*px* \sqcap - ?*pt2sy* \sqcap *y*)
 by *auto*
 have $p0[?px - p0^{T^*} * (?pyppy[[y]]) \mapsto ?p2t] = p0[?px - p0^{T^*} * (p[[?py]]) \mapsto ?p2t]$
 using 1 *put-get vector-mult-closed path-halving-invariant-def* by *auto*
 also have ... = $p0[?px - ?pt2sy \mapsto ?p2t]$
 using 1 *comp-associative path-halving-invariant-aux(2)* by *force*
 also have ... = $p0[?pxy \sqcup (?px \sqcap - ?pt2sy \sqcap ?pty) \sqcup (?px \sqcap - ?pt2sy \sqcap y) \mapsto ?p2t]$
 using 13 by (*metis comp-inf.semiring.distrib-left inf.sup-monoid.add-assoc*)
 also have ... = $(?p[?px \sqcap - ?pt2sy \sqcap ?pty \mapsto ?p2t])[?px \sqcap - ?pt2sy \sqcap y \mapsto ?p2t]$
 using 15 by (*smt (z3) update-same-3 comp-inf.semiring.mult-not-zero inf.sup-monoid.add-assoc inf.sup-monoid.add-commute*)
 also have ... = $(p[?px \sqcap - ?pt2sy \sqcap ?pty \mapsto ?p2t])[?px \sqcap - ?pt2sy \sqcap y \mapsto ?p2t]$
 using 1 *path-halving-invariant-def* by *auto*
 also have ... = $p[?px \sqcap - ?pt2sy \sqcap y \mapsto ?p2t]$
 using 6 by *simp*
 also have ... = ?*pyppy*
 using 9 by *auto*
 finally show ?*thesis*
 .
 qed
 show *univalent* *p0 total* *p0 acyclic* (*p0 - 1*)
 using 1 *path-halving-invariant-def* by *auto*
 qed
 let ?*s* = { *z* . *regular* *z* \wedge *z* \leq $p^{T^*} * y$ }
 let ?*t* = { *z* . *regular* *z* \wedge *z* \leq ?*pyppy* $^{T^*} * (?pyppy[[y]])$ }
 have ?*pyppy* $^{T^*} * (?pyppy[[y]]) = ?pyppy^{T^*} * (p[[?py]])$
 using 1 *put-get vector-mult-closed path-halving-invariant-def* by *force*
 also have ... $\leq p^{+T^*} * (p[[?py]])$
 using 1 *path-halving-invariant-def update-square-plus conv-order mult-left-isotone star-isotone* by *force*
 also have ... = $p^{T^*} * p^T * p^T * y$
 by (*simp add: conv-plus-commute star.left-plus-circ mult-assoc*)
 also have ... $\leq p^{T^+} * y$
 by (*metis mult-left-isotone star.left-plus-below-circ star-plus*)
 finally have 16: ?*pyppy* $^{T^*} * (?pyppy[[y]]) \leq p^{T^+} * y$
 .
 hence ?*pyppy* $^{T^*} * (?pyppy[[y]]) \leq p^{T^*} * y$
 using *mult-left-isotone order-lesseq-imp star.left-plus-below-circ* by *blast*

```

hence 17:  $?t \subseteq ?s$ 
  using order-trans by auto
have 18:  $y \in ?s$ 
  using 1 bijjective-regular path-compression-1b path-halving-invariant-def by
force
have 19:  $\neg y \in ?t$ 
proof
  assume  $y \in ?t$ 
  hence  $y \leq ?pyppy^{T*} * (?pyppy[[y]])$ 
    by simp
  hence  $y \leq p^{T+} * y$ 
    using 16 dual-order.trans by blast
  hence  $y = \text{root } p \ y$ 
    using 1 find-set-precondition-def loop-root-2 path-halving-invariant-def by
blast
  hence  $y = ?py$ 
    using 1 by (metis find-set-precondition-def root-successor-loop
path-halving-invariant-def)
  thus False
    using 1 by simp
qed
show  $\text{card } ?t < \text{card } ?s$ 
  apply (rule psubset-card-mono)
  subgoal using finite-regular by simp
  subgoal using 17 18 19 by auto
  done
qed
qed

lemma path-halving-3:
  path-halving-invariant  $p \ x \ y \ p0 \wedge y = p[[y]] \implies \text{path-halving-postcondition } p \ x \ y \ p0$ 
proof –
  assume 1: path-halving-invariant  $p \ x \ y \ p0 \wedge y = p[[y]]$ 
  show path-halving-postcondition  $p \ x \ y \ p0$ 
  proof (unfold path-halving-postcondition-def path-compression-precondition-def,
intro conjI)
    show univalent  $p$  total  $p$  acyclic  $(p - 1)$ 
      using 1 find-set-precondition-def path-halving-invariant-def by blast+
    show vector  $x$  injective  $x$  surjective  $x$ 
      using 1 find-set-precondition-def path-halving-invariant-def by blast+
    show 2: vector  $y$  injective  $y$  surjective  $y$ 
      using 1 path-halving-invariant-def by blast+
    have find-set-invariant  $p \ x \ y$ 
      using 1 find-set-invariant-def path-halving-invariant-def by blast
    thus  $y = \text{root } p \ x$ 
      using 1 find-set-3 find-set-postcondition-def by blast
    show  $p \sqcap 1 = p0 \sqcap 1$ 
      using 1 path-halving-invariant-aux(4) by blast

```



```

show fc p = fc p0
  using 1 path-halving-invariant-aux(5) by blast
have 3: y = p0[[y]]
  using 1 path-halving-invariant-aux(1) by auto
hence p0T* * y = y
  using order.antisym path-compression-1b star-left-induct-mult-equal by auto
hence 4: p0[(p0 * p0)T* * x - y → (p0 * p0)T] = p
  using 1 path-halving-invariant-def by auto
have (p0 * p0)T * y = y
  using 3 mult-assoc conv-dist-comp by auto
hence y ⊓ p0 * p0 = y ⊓ p0
  using 2 3 by (metis update-postcondition)
hence 5: y ⊓ p = y ⊓ p0 * p0
  using 1 2 3 by (smt update-postcondition)
have p0[(p0 * p0)T* * x → (p0 * p0)T] = (p0[(p0 * p0)T* * x - y → (p0 *
p0)T])[(p0 * p0)T* * x ⊓ y → (p0 * p0)T]
  using 1 bijective-regular path-halving-invariant-def update-split by blast
also have ... = p[(p0 * p0)T* * x ⊓ y → (p0 * p0)T]
  using 4 by simp
also have ... = p
  apply (rule update-same-sub)
  using 5 apply simp
  apply simp
  using 1 bijective-regular inf.absorb2 path-halving-invariant-def by auto
finally show p0[(p0 * p0)T* * x → (p0 * p0)T] = p
  .
qed
qed

```

```

theorem find-path-halving:
  VARS p y
  [ find-set-precondition p x ∧ p0 = p ]
  y := x;
  WHILE y ≠ p[[y]]
    INV { path-halving-invariant p x y p0 }
    VAR { (pT* * y)↓ }
    DO p[y] := p[[p[[y]]]];
    y := p[[y]]
  OD
  [ path-halving-postcondition p x y p0 ]
  apply vcg-tc-simp
  apply (fact path-halving-1)
  apply (fact path-halving-2)
  by (fact path-halving-3)

```

6.3 Path Splitting

Path splitting is another variant of the path compression technique. We implement it again independently of find-set, using a second while-loop which

iterates over the same path to the root. We prove that path splitting preserves the equivalence-relational semantics of the disjoint-set forest and also preserves the roots of the component trees. Additionally we prove the exact effect of path splitting, which is to replace every parent pointer with a pointer to the respective grandparent.

definition *path-splitting-invariant* $p\ x\ y\ p0 \equiv$
find-set-precondition $p\ x \wedge \text{point } y \wedge y \leq p0^{T^*} * x \wedge$
 $p0[p0^{T^*} * x - p0^{T^*} * y \mapsto (p0 * p0)^T] = p \wedge$
disjoint-set-forest $p0$

definition *path-splitting-postcondition* $p\ x\ y\ p0 \equiv$
path-compression-precondition $p\ x\ y \wedge p \sqcap 1 = p0 \sqcap 1 \wedge \text{fc } p = \text{fc } p0 \wedge$
 $p0[p0^{T^*} * x \mapsto (p0 * p0)^T] = p$

lemma *path-splitting-invariant-aux-1*:

assumes *point* x

and *point* y

and *disjoint-set-forest* $p0$

shows $(p0[p0^{T^*} * x - p0^{T^*} * y \mapsto (p0 * p0)^T]) \sqcap 1 = p0 \sqcap 1$

and $\text{fc } (p0[p0^{T^*} * x - p0^{T^*} * y \mapsto (p0 * p0)^T]) = \text{fc } p0$

and $p0^{T^*} * x \leq p0^* * \text{root } p0\ x$

proof –

let $?p2 = p0 * p0$

let $?p2t = ?p2^T$

let $?px = p0^{T^*} * x$

let $?py = -(p0^{T^*} * y)$

let $?pxy = ?px \sqcap ?py$

let $?q1 = ?pxy \sqcap p0$

let $?q2 = -?pxy \sqcap p0$

let $?q3 = ?pxy \sqcap ?p2$

let $?q4 = -?pxy \sqcap ?p2$

let $?p = p0[?pxy \mapsto ?p2t]$

let $?r0 = \text{root } p0\ x$

let $?rp = \text{root } ?p\ x$

have 1: *regular* $?px \wedge \text{regular } (p0^{T^*} * y) \wedge \text{regular } ?pxy$

using *assms* *bijective-regular* *find-set-precondition-def* *mapping-regular*

pp-dist-comp *regular-closed-star* *regular-conv-closed* *path-halving-invariant-def*

regular-closed-inf **by** *auto*

have 2: *vector* $x \wedge \text{vector } ?px \wedge \text{vector } ?py \wedge \text{vector } ?pxy$

using *assms*(1,2) *find-set-precondition-def* *vector-complement-closed*

vector-mult-closed *path-halving-invariant-def* *vector-inf-closed* **by** *auto*

have 3: $?r0 \leq p0 * ?r0$

by (*metis* *assms*(3) *dedekind-1* *inf.le-iff-sup* *root-successor-loop* *top-greatest*)

hence $?pxy \sqcap p0 * ?r0 \leq ?pxy \sqcap ?p2 * ?r0$

by (*metis* *comp-associative* *inf.eq-refl* *inf.sup-right-isotone* *mult-isotone*)

hence 4: $?q1 * ?r0 \leq ?q3 * ?r0$

using 2 **by** (*simp* *add*: *vector-inf-comp*)

have 5: $?q1 * ?q2 \leq ?q3$

using 2 **by** (*smt* (z3) *comp-isotone* *inf.cobounded1* *inf.cobounded2* *inf-greatest*)

vector-export-comp)
have $?q1 * ?q2^* * ?r0 = ?q1 * ?r0 \sqcup ?q1 * ?q2 * ?q2^* * ?r0$
by (*metis comp-associative semiring.distrib-left star.circ-loop-fixpoint sup-commute*)
also have $\dots \leq ?q1 * ?r0 \sqcup ?q3 * ?q2^* * ?r0$
using 5 **by** (*meson mult-left-isotone sup-right-isotone*)
also have $\dots \leq ?q3 * ?r0 \sqcup ?q3 * ?q2^* * ?r0$
using 4 *sup-left-isotone* **by** *blast*
also have $\dots = ?q3 * ?q2^* * ?r0$
by (*smt (verit, del-insts) comp-associative semiring.distrib-left star.circ-loop-fixpoint star.circ-transitive-equal star-involutive sup-commute*)
finally have 6: $?q1 * ?q2^* * ?r0 \leq ?q3 * ?q2^* * ?r0$
.

have $?q1 * (-?pxy \sqcap p0^+) * ?pxy \leq (?px \sqcap p0) * (-?pxy \sqcap p0^+) * ?pxy$
by (*meson comp-inf.comp-left-subdist-inf inf.boundedE mult-left-isotone*)
also have $\dots \leq (?px \sqcap p0) * (-?pxy \sqcap p0^+) * ?py$
by (*simp add: mult-right-isotone*)
also have $\dots \leq ?px^T * (-?pxy \sqcap p0^+) * ?py$
proof –
have $?px \sqcap p0 \leq ?px^T * p0$
using 2 **by** (*simp add: vector-restrict-comp-conv*)
also have $\dots \leq ?px^T$
by (*metis comp-associative conv-dist-comp conv-involutive conv-star-commute mult-right-isotone star.circ-increasing star.circ-transitive-equal*)
finally show *?thesis*
using *mult-left-isotone* **by** *auto*
qed

also have $\dots = top * (?px \sqcap -?pxy \sqcap p0^+) * ?py$
using 2 **by** (*smt (z3) comp-inf.star-plus conv-dist-inf covector-inf-comp-3 inf-top.right-neutral vector-complement-closed vector-inf-closed*)
also have $\dots \leq top * (-?py \sqcap p0^+) * ?py$
by (*metis comp-inf.comp-isotone comp-isotone inf.cobounded2 inf.eq-refl inf-import-p*)
also have $\dots = top * (-?py \sqcap p0^+ \sqcap ?py^T) * top$
using 2 **by** (*simp add: comp-associative covector-inf-comp-3*)
also have $\dots = bot$
proof –
have $p0^{T^*} * y - y^T * p0^* = p0^{T^*} * y * y^T * -p0^*$
using 2 **by** (*metis assms(2) bijective-conv-mapping comp-mapping-complement vector-covector vector-export-comp vector-mult-closed*)
also have $\dots \leq p0^{T^*} * -p0^*$
by (*meson assms(2) mult-left-isotone order-refl shunt-bijective*)
also have $\dots \leq -p0^*$
by (*simp add: conv-complement conv-star-commute pp-increasing schroeder-6-p star.circ-transitive-equal*)
also have $\dots \leq -p0^+$
by (*simp add: p-antitone star.left-plus-below-circ*)
finally have $-?py \sqcap p0^+ \sqcap ?py^T = bot$
by (*metis comp-inf.p-pp-comp conv-complement conv-dist-comp*)

*conv-involutive conv-star-commute p-shunting-swap pp-isotone
pseudo-complement-pp regular-closed-p)*
thus *?thesis*
by *simp*
qed
finally have 7: $?q1 * (-?pxy \sqcap p0^+) * ?pxy = bot$
using *le-bot by blast*
have $?q2^+ \leq -?pxy$
using 2 **by** (*smt (z3) comp-isotone complement-conv-sub inf.order-trans
inf.sup-right-divisibility inf-commute symmetric-top-closed top-greatest*)
hence $?q2^+ \leq -?pxy \sqcap p0^+$
by (*simp add: comp-isotone star-isotone*)
hence 8: $?q1 * ?q2^+ * ?pxy = bot$
using 7 *mult-left-isotone mult-right-isotone le-bot* **by** *auto*
have $?q1 * ?q2^+ * ?q3^* = ?q1 * ?q2^+ \sqcup ?q1 * ?q2^+ * ?q3^+$
by (*smt (z3) comp-associative star.circ-back-loop-fixpoint star.circ-plus-same
sup-commute*)
also have $\dots \leq ?q1 * ?q2^+ \sqcup ?q1 * ?q2^+ * ?pxy$
using 2 **by** (*smt (z3) inf.cobounded1 mult-right-isotone sup-right-isotone
vector-inf-comp*)
finally have 9: $?q1 * ?q2^+ * ?q3^* \leq ?q1 * ?q2^+$
using 8 **by** *simp*
have 10: $?q1 * ?q4 * ?pxy = bot$
proof –
have $?p2 \leq p0^+$
by (*simp add: mult-right-isotone star.circ-increasing*)
thus *?thesis*
using 7 **by** (*metis mult-left-isotone mult-right-isotone le-bot
comp-inf.comp-isotone eq-refl*)
qed
have 11: $?q1 * ?q2 * ?pxy = bot$
proof –
have $p0 \leq p0^+$
by (*simp add: star.circ-mult-increasing*)
thus *?thesis*
using 7 **by** (*metis mult-left-isotone mult-right-isotone le-bot
comp-inf.comp-isotone eq-refl*)
qed
have 12: $?q2 \leq p0 * ?q3^* * ?q2^*$
by (*smt (verit, del-insts) conv-dist-comp conv-order conv-star-commute
inf.coboundedI1 inf.orderE inf.sup-monoid.add-commute path-compression-1b*)
have $?q3 * p0 * ?q3^* * ?q2^* = ?q1 * p0 * p0 * ?q3^* * ?q2^*$
using 2 *vector-inf-comp* **by** *auto*
also have $\dots = ?q1 * (?q3 \sqcup ?q4) * ?q3^* * ?q2^*$
using 1 **by** (*smt (z3) comp-associative comp-inf.mult-right-dist-sup
comp-inf.star-slide inf-top.right-neutral regular-complement-top*)
also have $\dots = ?q1 * ?q3 * ?q3^* * ?q2^* \sqcup ?q1 * ?q4 * ?q3^* * ?q2^*$
using *mult-left-dist-sup mult-right-dist-sup* **by** *auto*
also have $\dots \leq ?q1 * ?q3^* * ?q2^* \sqcup ?q1 * ?q4 * ?q3^* * ?q2^*$

by (*smt (z3) mult-left-isotone mult-left-sub-dist-sup-right sup-left-isotone sup-right-divisibility mult-assoc star.left-plus-below-circ*)
also have ... = ?q1 * ?q3* * ?q2* \sqcup ?q1 * ?q4 * ?q2* \sqcup ?q1 * ?q4 * ?q3+ * ?q2*
by (*smt (z3) semiring.combine-common-factor star.circ-back-loop-fixpoint star-plus sup-monoid.add-commute mult-assoc*)
also have ... \leq ?q1 * ?q3* * ?q2* \sqcup ?q1 * ?q4 * ?q2* \sqcup ?q1 * ?q4 * ?pxy * ?q3* * ?q2*
by (*smt (verit, ccfv-threshold) comp-isotone inf.sup-right-divisibility inf-commute order.refl semiring.add-left-mono mult-assoc*)
also have ... = ?q1 * ?q3* * ?q2* \sqcup ?q1 * ?q4 * ?q2*
using 10 by simp
also have ... = ?q1 * ?q3* * ?q2* \sqcup ?q1 * ?q2 * p0 * ?q2*
using 2 by (*smt vector-complement-closed vector-inf-comp mult-assoc*)
also have ... = ?q1 * ?q3* * ?q2* \sqcup ?q1 * ?q2 * (?q2 \sqcup ?q1) * ?q2*
using 1 by (*smt (z3) comp-associative comp-inf.mult-right-dist-sup comp-inf.star-slide inf-top.right-neutral regular-complement-top*)
also have ... = ?q1 * ?q3* * ?q2* \sqcup ?q1 * ?q2 * ?q2 * ?q2* \sqcup ?q1 * ?q2 * ?q1 * ?q2*
using mult-left-dist-sup mult-right-dist-sup sup-commute sup-left-commute by auto
also have ... \leq ?q1 * ?q3* * ?q2* \sqcup ?q1 * ?q2 * ?q2 * ?q2* \sqcup ?q1 * ?q2 * ?pxy * ?q2*
by (*smt (verit, ccfv-threshold) comp-isotone inf.sup-right-divisibility inf-commute order.refl semiring.add-left-mono mult-assoc*)
also have ... = ?q1 * ?q3* * ?q2* \sqcup ?q1 * ?q2 * ?q2 * ?q2*
using 11 by simp
also have ... \leq ?q1 * ?q3* * ?q2* \sqcup ?q1 * ?q2*
by (*smt comp-associative comp-isotone mult-right-isotone star.circ-increasing star.circ-transitive-equal star.left-plus-below-circ sup-right-isotone*)
also have ... = ?q1 * ?q3* * ?q2*
by (*smt (verit, best) comp-associative semiring.distrib-left star.circ-loop-fixpoint star.circ-transitive-equal star-involutive*)
finally have 13: ?q3 * p0 * ?q3* * ?q2* \leq p0 * ?q3* * ?q2*
by (*meson inf.cobounded2 mult-left-isotone order-lesseq-imp*)
hence ?q3 * p0 * ?q3* * ?q2* \sqcup ?q2 \leq p0 * ?q3* * ?q2*
using 12 by simp
hence ?q3* * ?q2 \leq p0 * ?q3* * ?q2*
by (*simp add: star-left-induct mult-assoc*)
hence ?q1 * ?q3* * ?q2 \leq ?q1 * p0 * ?q3* * ?q2*
by (*simp add: comp-associative mult-right-isotone*)
hence ?q1 * ?q3* * ?q2 \leq ?q3+ * ?q2*
using 2 by (*simp add: vector-inf-comp*)
hence 14: ?q1 * ?q3* * ?q2 \leq ?q3* * ?q2*
using mult-left-isotone order-lesseq-imp star.left-plus-below-circ by blast
have p0 * ?r0 \leq p0 * ?q3* * ?q2* * ?r0
by (*metis comp-associative mult-1-right mult-left-isotone mult-right-isotone reflexive-mult-closed star.circ-reflexive*)
hence 15: ?r0 \leq p0 * ?q3* * ?q2* * ?r0

using 3 dual-order.trans by blast
 have $?q3 * p0 * ?q3^* * ?q2^* * ?r0 \leq p0 * ?q3^* * ?q2^* * ?r0$
 using 13 mult-left-isotone by blast
 hence $?q3 * p0 * ?q3^* * ?q2^* * ?r0 \sqcup ?r0 \leq p0 * ?q3^* * ?q2^* * ?r0$
 using 15 by simp
 hence $?q3^* * ?r0 \leq p0 * ?q3^* * ?q2^* * ?r0$
 by (simp add: star-left-induct mult-assoc)
 hence $?q1 * ?q3^* * ?r0 \leq ?q1 * p0 * ?q3^* * ?q2^* * ?r0$
 by (simp add: comp-associative mult-right-isotone)
 hence $?q1 * ?q3^* * ?r0 \leq ?q3^+ * ?q2^* * ?r0$
 using 2 by (simp add: vector-inf-comp)
 hence 16: $?q1 * ?q3^* * ?r0 \leq ?q3^* * ?q2^* * ?r0$
 using mult-left-isotone order-lesseq-imp star.left-plus-below-circ by blast
 have $?q1 * ?q3^* * ?q2^* * ?r0 = ?q1 * ?q3^* * ?r0 \sqcup ?q1 * ?q3^* * ?q2^+ * ?r0$
 by (smt (z3) comp-associative mult-right-dist-sup star.circ-back-loop-fixpoint
 star.circ-plus-same sup-commute)
 also have $\dots \leq ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q3^* * ?q2^+ * ?r0$
 using 16 sup-left-isotone by blast
 also have $\dots \leq ?q3^* * ?q2^* * ?r0 \sqcup ?q3^* * ?q2^* * ?q2^* * ?r0$
 using 14 by (smt (z3) inf.eq-refl semiring.distrib-right
 star.circ-transitive-equal sup.absorb2 sup-monoid.add-commute mult-assoc)
 also have $\dots = ?q3^* * ?q2^* * ?r0$
 by (simp add: comp-associative star.circ-transitive-equal)
 finally have 17: $?q1 * ?q3^* * ?q2^* * ?r0 \leq ?q3^* * ?q2^* * ?r0$
 .
 have $?r0 \leq ?q2^* * ?r0$
 using star.circ-loop-fixpoint sup-right-divisibility by auto
 also have $\dots \leq ?q3^* * ?q2^* * ?r0$
 using comp-associative star.circ-loop-fixpoint sup-right-divisibility by force
 also have $\dots \leq ?q2^* * ?q3^* * ?q2^* * ?r0$
 using comp-associative star.circ-loop-fixpoint sup-right-divisibility by force
 finally have 18: $?r0 \leq ?q2^* * ?q3^* * ?q2^* * ?r0$
 .
 have $p0 * ?q2^* * ?q3^* * ?q2^* * ?r0 = (?q2 \sqcup ?q1) * ?q2^* * ?q3^* * ?q2^* * ?r0$
 using 1 by (smt (z3) comp-inf.mult-right-dist-sup comp-inf.star-plus
 inf-top.right-neutral regular-complement-top)
 also have $\dots = ?q2 * ?q2^* * ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q2^* * ?q3^* * ?q2^* * ?r0$
 using mult-right-dist-sup by auto
 also have $\dots \leq ?q2^* * ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q2^* * ?q3^* * ?q2^* * ?r0$
 by (smt (z3) comp-left-increasing-sup star.circ-loop-fixpoint sup-left-isotone
 mult-assoc)
 also have $\dots = ?q2^* * ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q2^+ * ?q3^* * ?q2^* * ?r0$
 by (smt (z3) mult-left-dist-sup semiring.combine-common-factor
 star.circ-loop-fixpoint sup-monoid.add-commute mult-assoc)
 also have $\dots \leq ?q2^* * ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q2^+ * ?q3^* * ?q2^* * ?r0$
 using 9 mult-left-isotone sup-right-isotone by auto

also have ... $\leq ?q2^* * ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q2^* * ?r0$
by (smt (z3) comp-associative comp-isotone inf.eq-refl semiring.add-right-mono star.circ-transitive-equal star.left-plus-below-circ sup-commute)
also have ... $\leq ?q2^* * ?q3^* * ?q2^* * ?r0 \sqcup ?q1 * ?q3^* * ?q2^* * ?r0 \sqcup ?q3 * ?q2^* * ?r0$
using 6 sup-right-isotone **by** blast
also have ... $= ?q2^* * ?q3^* * ?q2^* * ?r0 \sqcup ?q3 * ?q2^* * ?r0$
using 17 **by** (smt (z3) le-iff-sup semiring.combine-common-factor semiring.distrib-right star.circ-loop-fixpoint sup-monoid.add-commute)
also have ... $\leq ?q2^* * ?q3^* * ?q2^* * ?r0 \sqcup ?q3^* * ?q2^* * ?r0$
by (meson mult-left-isotone star.circ-increasing sup-right-isotone)
also have ... $= ?q2^* * ?q3^* * ?q2^* * ?r0$
by (smt (z3) comp-associative star.circ-loop-fixpoint star.circ-transitive-equal star-involutive)
finally have $p0 * ?q2^* * ?q3^* * ?q2^* * ?r0 \sqcup ?r0 \leq ?q2^* * ?q3^* * ?q2^* * ?r0$
using 18 sup.boundedI **by** blast
hence $p0^* * ?r0 \leq ?q2^* * ?q3^* * ?q2^* * ?r0$
by (simp add: comp-associative star-left-induct)
also have ... $\leq ?p^* * ?q3^* * ?q2^* * ?r0$
by (metis mult-left-isotone star.circ-sub-dist sup-commute)
also have ... $\leq ?p^* * ?p^* * ?q2^* * ?r0$
by (simp add: mult-left-isotone mult-right-isotone star-isotone)
also have ... $\leq ?p^* * ?p^* * ?p^* * ?r0$
by (metis mult-isotone order.refl star.circ-sub-dist sup-commute)
finally have 19: $p0^* * ?r0 \leq ?p^* * ?r0$
by (simp add: star.circ-transitive-equal)
have 20: $?p^* \leq p0^*$
by (metis star.left-plus-circ star-isotone update-square-ub-plus)
hence 21: $p0^* * ?r0 = ?p^* * ?r0$
using 19 order.antisym mult-left-isotone **by** auto
have $?p \sqcap 1 = (?q3 \sqcap 1) \sqcup (?q2 \sqcap 1)$
using comp-inf.semiring.distrib-right conv-involutive **by** auto
also have ... $= (?q1 \sqcap 1) \sqcup (?q2 \sqcap 1)$
using assms(3) acyclic-square path-splitting-invariant-def
inf.sup-monoid.add-assoc **by** auto
also have ... $= (?pxy \sqcup -?pxy) \sqcap p0 \sqcap 1$
using inf-sup-distrib2 **by** auto
also have ... $= p0 \sqcap 1$
using 1 **by** (metis inf.sup-monoid.add-commute inf-sup-distrib1 maddux-3-11-pp)
finally show 22: $?p \sqcap 1 = p0 \sqcap 1$
.
have $?p^{T^*} * x \leq p0^{T^*} * x$
using 20 **by** (metis conv-isotone conv-star-commute mult-left-isotone)
hence 23: $?rp \leq ?r0$
using 22 comp-inf.mult-left-isotone **by** auto
have 24: disjoint-set-forest ?p

using 1 2 *assms(3) disjoint-set-forest-update-square* **by** *blast*
hence 25: *point ?rp*
using *root-point assms(1)* **by** *auto*
have $?r0 * ?rp^T = ?r0 * x^T * ?p^* * (?p \sqcap 1)$
by (*smt (z3) comp-associative conv-dist-comp conv-dist-inf conv-involutive conv-star-commute inf.sup-monoid.add-commute one-inf-conv root-var star-one star-sup-one wcc-one*)
also have $\dots \leq (p0 \sqcap 1) * p0^{T^*} * 1 * ?p^* * (?p \sqcap 1)$
by (*smt (z3) assms(1) comp-associative mult-left-isotone mult-right-isotone root-var*)
also have $\dots \leq (p0 \sqcap 1) * p0^{T^*} * p0^* * (p0 \sqcap 1)$
using 20 22 *comp-isotone* **by** *force*
also have $\dots = (p0 \sqcap 1) * p0^* * (p0 \sqcap 1) \sqcup (p0 \sqcap 1) * p0^{T^*} * (p0 \sqcap 1)$
by (*simp add: assms(3) cancel-separate-eq sup-monoid.add-commute mult-assoc mult-left-dist-sup semiring.distrib-right*)
also have $\dots = (p0 \sqcap 1) * (p0 \sqcap 1) \sqcup (p0 \sqcap 1) * p0^{T^*} * (p0 \sqcap 1)$
using *univalent-root-successors assms(3)* **by** *simp*
also have $\dots = (p0 \sqcap 1) * (p0 \sqcap 1) \sqcup (p0 \sqcap 1) * ((p0 \sqcap 1) * p0^*)^T$
by (*smt (z3) comp-associative conv-dist-comp conv-dist-inf conv-star-commute inf.sup-monoid.add-commute one-inf-conv star-one star-sup-one wcc-one*)
also have $\dots = (p0 \sqcap 1) * (p0 \sqcap 1)$
by (*metis univalent-root-successors assms(3) conv-dist-inf inf.sup-monoid.add-commute one-inf-conv sup-idem symmetric-one-closed*)
also have $\dots \leq 1$
by (*simp add: coreflexive-mult-closed*)
finally have $?r0 * ?rp^T \leq 1$
.

hence $?r0 \leq 1 * ?rp$
using 25 *shunt-bijective* **by** *blast*
hence 26: $?r0 = ?rp$
using 23 *order.antisym* **by** *simp*
have $?px * ?r0^T = ?px * x^T * p0^* * (p0 \sqcap 1)$
by (*smt (z3) comp-associative conv-dist-comp conv-dist-inf conv-involutive conv-star-commute inf.sup-monoid.add-commute one-inf-conv root-var star-one star-sup-one wcc-one*)
also have $\dots \leq p0^{T^*} * 1 * p0^* * (p0 \sqcap 1)$
by (*smt (z3) assms(1) comp-associative mult-left-isotone mult-right-isotone root-var*)
also have $\dots = p0^* * (p0 \sqcap 1) \sqcup p0^{T^*} * (p0 \sqcap 1)$
by (*simp add: assms(3) cancel-separate-eq sup-monoid.add-commute mult-right-dist-sup*)
also have $\dots = p0^* * (p0 \sqcap 1) \sqcup ((p0 \sqcap 1) * p0^*)^T$
by (*smt (z3) conv-dist-comp conv-dist-inf conv-star-commute inf.sup-monoid.add-commute one-inf-conv star-one star-sup-one wcc-one*)
also have $\dots = p0^* * (p0 \sqcap 1) \sqcup (p0 \sqcap 1)$
by (*metis univalent-root-successors assms(3) conv-dist-inf inf.sup-monoid.add-commute one-inf-conv symmetric-one-closed*)
also have $\dots = p0^* * (p0 \sqcap 1)$
by (*metis conv-involutive path-compression-1b sup.absorb2 sup-commute*)

also have $\dots \leq p0^*$
by (*simp add: inf.coboundedI1 star.circ-increasing star.circ-mult-upper-bound*)
finally have 27: $?px * ?r0^T \leq p0^*$
 .
thus 28: $?px \leq p0^* * ?r0$
by (*simp add: assms(1,3) root-point shunt-bijective*)
have 29: *point ?r0*
using *root-point assms(1,3) by auto*
hence 30: *mapping (?r0^T)*
using *bijective-conv-mapping by blast*
have $?r0 * (?px \sqcap p0) = ?r0 * top * (?px \sqcap p0)$
using 29 *by force*
also have $\dots = ?r0 * ?px^T * p0$
using 29 *by (metis assms(1) covector-inf-comp-3 vector-covector vector-mult-closed)*
also have $\dots = ?r0 * x^T * p0^* * p0$
using *comp-associative conv-dist-comp conv-star-commute by auto*
also have $\dots \leq ?r0 * x^T * p0^*$
by (*simp add: comp-associative mult-right-isotone star.circ-plus-same star.left-plus-below-circ*)
also have $\dots = ?r0 * ?px^T$
by (*simp add: comp-associative conv-dist-comp conv-star-commute*)
also have $\dots = (?px * ?r0^T)^T$
by (*simp add: conv-dist-comp*)
also have $\dots \leq p0^{T*}$
using 27 *conv-isotone conv-star-commute by fastforce*
finally have $?r0 * (?px \sqcap p0) \leq p0^{T*}$
 .
hence $?px \sqcap p0 \leq ?r0^T * p0^{T*}$
using 30 *shunt-mapping by auto*
hence $?px \sqcap p0 \leq p0^* * ?r0 \sqcap ?r0^T * p0^{T*}$
using 28 *inf.coboundedI2 inf.sup-monoid.add-commute by fastforce*
also have $\dots = p0^* * ?r0 * ?r0^T * p0^{T*}$
using 29 *by (smt (z3) vector-covector vector-inf-comp vector-mult-closed)*
also have $\dots = ?p^* * ?r0 * ?r0^T * ?p^{T*}$
using 21 *by (smt comp-associative conv-dist-comp conv-star-commute)*
also have $\dots = ?p^* * ?rp * ?rp^T * ?p^{T*}$
using 26 *by auto*
also have $\dots \leq ?p^* * 1 * ?p^{T*}$
using 25 *by (smt (z3) comp-associative mult-left-isotone mult-right-isotone)*
finally have 31: $?px \sqcap p0 \leq fc ?p$
by auto
have $-?px \sqcap p0 \leq ?p$
by (*simp add: inf.sup-monoid.add-commute le-supI1 sup-commute*)
also have $\dots \leq fc ?p$
using *fc-increasing by auto*
finally have $p0 \leq fc ?p$
using 1 31 *by (smt (z3) inf.sup-monoid.add-commute maddux-3-11-pp semiring.add-left-mono sup.orderE sup-commute)*

also have $\dots \leq wcc \ ?p$
using *star.circ-sub-dist-3* **by** *auto*
finally have $32: wcc \ p0 \leq wcc \ ?p$
using *wcc-below-wcc* **by** *blast*
have $?p \leq wcc \ p0$
by (*simp add: inf.coboundedI1 inf.sup-monoid.add-commute*
star.circ-mult-upper-bound star.circ-sub-dist-1)
hence $wcc \ ?p \leq wcc \ p0$
using *wcc-below-wcc* **by** *blast*
hence $wcc \ ?p = wcc \ p0$
using 32 *order.antisym* **by** *blast*
thus $fc \ ?p = fc \ p0$
using 24 *assms(3)* *fc-wcc* **by** *auto*
qed

lemma *path-splitting-invariant-aux:*

assumes *path-splitting-invariant* $p \ x \ y \ p0$
shows $p[[y]] = p0[[y]]$
and $p[[p[[y]]]] = p0[[p0[[y]]]]$
and $p[[p[[p[[y]]]]]] = p0[[p0[[p0[[y]]]]]]$
and $p \sqcap 1 = p0 \sqcap 1$
and $fc \ p = fc \ p0$
proof –
let $?p2 = p0 * p0$
let $?p2t = ?p2^T$
let $?px = p0^{T*} * x$
let $?py = -(p0^{T*} * y)$
let $?pxy = ?px \sqcap ?py$
let $?p = p0[?pxy \mapsto ?p2t]$
have $?p[[y]] = p0[[y]]$
apply (*rule put-get-different-vector*)
using *assms find-set-precondition-def vector-complement-closed*
vector-inf-closed vector-mult-closed path-splitting-invariant-def **apply** *force*
by (*meson inf.cobounded2 order-lesseq-imp p-antitone-iff path-compression-1b*)
thus $1: p[[y]] = p0[[y]]$
using *assms path-splitting-invariant-def* **by** *auto*
have $?p[[p0[[y]]]] = p0[[p0[[y]]]]$
apply (*rule put-get-different-vector*)
using *assms find-set-precondition-def vector-complement-closed*
vector-inf-closed vector-mult-closed path-splitting-invariant-def **apply** *force*
by (*metis comp-isotone inf.boundedE inf.coboundedI2 inf.eq-refl p-antitone-iff*
selection-closed-id star.circ-increasing)
thus $2: p[[p[[y]]]] = p0[[p0[[y]]]]$
using 1 *assms path-splitting-invariant-def* **by** *auto*
have $?p[[p0[[p0[[y]]]]]] = p0[[p0[[p0[[y]]]]]]$
apply (*rule put-get-different-vector*)
using *assms find-set-precondition-def vector-complement-closed*
vector-inf-closed vector-mult-closed path-splitting-invariant-def **apply** *force*
by (*metis comp-associative comp-isotone conv-dist-comp conv-involutive*)

```

conv-order inf.coboundedI2 inf.le-iff-sup mult-left-isotone p-antitone-iff
p-antitone-inf star.circ-increasing star.circ-transitive-equal)
  thus p[[p[[p[[[y]]]]]]] = p0[[p0[[p0[[[y]]]]]]]
    using 2 assms path-splitting-invariant-def by auto
  show p  $\sqcap$  1 = p0  $\sqcap$  1
    using assms path-splitting-invariant-aux-1(1) path-splitting-invariant-def
find-set-precondition-def by auto
  show fc p = fc p0
    using assms path-splitting-invariant-aux-1(2) path-splitting-invariant-def
find-set-precondition-def by auto
qed

```

lemma *path-splitting-1:*

```

find-set-precondition p0 x  $\implies$  path-splitting-invariant p0 x x p0
proof –
  assume 1: find-set-precondition p0 x
  show path-splitting-invariant p0 x x p0
  proof (unfold path-splitting-invariant-def, intro conjI)
    show find-set-precondition p0 x
      using 1 by simp
    show vector x injective x surjective x
      using 1 find-set-precondition-def by auto
    show x  $\leq$  p0T* * x
      by (simp add: path-compression-1b)
    have (p0 * p0)T* * x  $\leq$  p0T* * x
      by (simp add: conv-dist-comp mult-left-isotone star.circ-square)
    thus p0[p0T* * x – p0T* * x  $\longrightarrow$  (p0 * p0)T] = p0
      by (smt (z3) inf.le-iff-sup inf-commute maddux-3-11-pp p-antitone-inf
pseudo-complement)
    show univalent p0 total p0 acyclic (p0 – 1)
      using 1 find-set-precondition-def by auto
  qed
qed

```

lemma *path-splitting-2:*

```

path-splitting-invariant p x y p0  $\wedge$  y  $\neq$  p[[y]]  $\implies$  path-splitting-invariant
(p[y $\longmapsto$ p[[p[[[y]]]]]]) x (p[[y]]) p0  $\wedge$  ((p[y $\longmapsto$ p[[p[[[y]]]]]])T* * (p[[y]])) $\downarrow$  < (pT* * y) $\downarrow$ 
proof –
  let ?py = p[[y]]
  let ?ppy = p[[?py]]
  let ?pyppy = p[y $\longmapsto$ ?ppy]
  let ?p2 = p0 * p0
  let ?p2t = ?p2T
  let ?p2ts = ?p2t*
  let ?px = p0T* * x
  let ?py2 = –(p0T* * y)
  let ?pxy = ?px  $\sqcap$  ?py2
  let ?p = p0[?pxy $\longrightarrow$ ?p2t]
  let ?pty = p0T * y

```

```

let ?pt2y = p0T * p0T * y
let ?pt2sy = p0T* * p0T * p0T * y
let ?ptpy = p0T+ * y
assume 1: path-splitting-invariant p x y p0 ∧ y ≠ ?py
have 2: point ?pty ∧ point ?pt2y
  using 1 by (smt (verit) comp-associative read-injective read-surjective
path-splitting-invariant-def)
show path-splitting-invariant ?pyppy x (p[[y]]) p0 ∧ (?pyppyT* * (p[[y]]))↓ <
(pT* * y)↓
proof
  show path-splitting-invariant ?pyppy x (p[[y]]) p0
  proof (unfold path-splitting-invariant-def, intro conjI)
    show 3: find-set-precondition ?pyppy x
    proof (unfold find-set-precondition-def, intro conjI)
      show univalent ?pyppy
      using 1 find-set-precondition-def read-injective update-univalent
path-splitting-invariant-def by auto
      show total ?pyppy
      using 1 bijective-regular find-set-precondition-def read-surjective
update-total path-splitting-invariant-def by force
      show acyclic (?pyppy - 1)
      apply (rule update-acyclic-3)
      using 1 find-set-precondition-def path-splitting-invariant-def apply blast
      using 1 2 comp-associative path-splitting-invariant-aux(2) apply force
      using 1 path-splitting-invariant-def apply blast
      by (metis inf.order-lesseq-imp mult-isotone star.circ-increasing
star.circ-square mult-assoc)
      show vector x injective x surjective x
      using 1 find-set-precondition-def path-splitting-invariant-def by auto
qed
show vector (p[[y]])
  using 1 comp-associative path-splitting-invariant-def by auto
show injective (p[[y]])
  using 1 3 read-injective path-splitting-invariant-def find-set-precondition-def
by auto
show surjective (p[[y]])
  using 1 3 read-surjective path-splitting-invariant-def
find-set-precondition-def by auto
show p[[y]] ≤ ?px
proof -
  have p[[y]] = p0[[y]]
    using 1 path-splitting-invariant-aux(1) by blast
  also have ... ≤ p0T * ?px
    using 1 path-splitting-invariant-def mult-right-isotone by force
  also have ... ≤ ?px
    by (metis comp-associative mult-left-isotone star.left-plus-below-circ)
  finally show ?thesis
  .
qed

```

```

show  $p0[?px - p0^{T^*} * (p[[y]]) \dashv\rightarrow ?p2t] = ?pyppy$ 
proof -
  have 4:  $p[?px \sqcap -?ptpy \sqcap y \dashv\rightarrow ?p2t] = ?pyppy$ 
  proof (cases  $y \leq -?ptpy$ )
    case True
      hence  $p[?px \sqcap -?ptpy \sqcap y \dashv\rightarrow ?p2t] = p[y \dashv\rightarrow ?p2t]$ 
      using 1 inf.absorb2 path-splitting-invariant-def by auto
      also have ... = ?pyppy
      using 1 by (metis comp-associative conv-dist-comp
        path-splitting-invariant-aux(2) path-splitting-invariant-def update-point-get)
      finally show ?thesis
    .
  next
    case False
      have vector ( $-?ptpy$ )
      using 1 vector-complement-closed vector-mult-closed
        path-splitting-invariant-def by blast
      hence 5:  $y \leq ?ptpy$ 
      using 1 by (smt (verit, del-insts) False bijective-regular
        point-in-vector-or-complement regular-closed-star regular-mult-closed
        total-conv-surjective univalent-conv-injective path-splitting-invariant-def)
      hence  $?px \sqcap -?ptpy \sqcap y = \text{bot}$ 
      by (simp add: inf.coboundedI2 p-antitone pseudo-complement)
      hence 6:  $p[?px \sqcap -?ptpy \sqcap y \dashv\rightarrow ?p2t] = p$ 
      by simp
      have 7:  $y = \text{root } p0 \ y$ 
      using 1 5 loop-root-2 path-splitting-invariant-def by blast
      have  $?pyppy = p[y \dashv\rightarrow p0[[p0[[y]]]]]$ 
      using 1 path-splitting-invariant-aux(2) by force
      also have ... =  $p[y \dashv\rightarrow p0[[y]]]$ 
      using 1 7 by (metis root-successor-loop path-splitting-invariant-def)
      also have ... =  $p[y \dashv\rightarrow ?py]$ 
      using 1 path-splitting-invariant-aux(1) by force
      also have ... =  $p$ 
      using 1 get-put path-splitting-invariant-def by blast
      finally show ?thesis
      using 6 by simp
  qed
  have 8:  $-?ptpy = ?py2 \sqcup (-?ptpy \sqcap y)$ 
  proof (rule order.antisym)
    have 9: regular ( $p0^{T^*} * y$ )  $\wedge$  regular ?ptpy
    using 1 bijective-regular mapping-conv-bijective pp-dist-star
      regular-mult-closed path-splitting-invariant-def by auto
    have  $p0^{T^*} * y \leq ?ptpy \sqcup y$ 
    by (simp add: star.circ-loop-fixpoint mult-assoc)
    also have ... =  $?ptpy \sqcup (-?ptpy \sqcap y)$ 
    using 9 by (metis maddux-3-21-pp)
    hence  $p0^{T^*} * y \sqcap -?ptpy \leq -?ptpy \sqcap y$ 
    using calculation half-shunting sup-commute by auto

```

thus $-?ptpy \leq ?py2 \sqcup (-?ptpy \sqcap y)$
using 9 by (*smt (z3) inf.sup-monoid.add-commute shunting-var-p*
sup.left-commute sup-commute)
have $-(p0^{T^*} * y) \leq -?ptpy$
by (*simp add: comp-isotone p-antitone star.left-plus-below-circ*)
thus $-(p0^{T^*} * y) \sqcup (-?ptpy \sqcap y) \leq -?ptpy$
by *simp*
qed
have *regular ?px regular y*
using 1 *bijjective-regular find-set-precondition-def mapping-regular*
pp-dist-comp regular-closed-star regular-conv-closed path-splitting-invariant-def **by**
auto
hence 10: *regular (?px \sqcap -?ptpy \sqcap y)*
by *auto*
have $p0[?px \sqcap -(p0^{T^*} * (p[[y]])) \mapsto ?p2t] = p0[?px \sqcap -?ptpy \mapsto ?p2t]$
using 1 by (*smt comp-associative path-splitting-invariant-aux(1)*)
star-plus
also have $\dots = p0[?pxy \sqcup (?px \sqcap -?ptpy \sqcap y) \mapsto ?p2t]$
using 8 by (*metis comp-inf.semiring.distrib-left*
inf.sup-monoid.add-assoc)
also have $\dots = ?p[?px \sqcap -?ptpy \sqcap y \mapsto ?p2t]$
using 10 by (*smt (z3) update-same comp-inf.semiring.mult-not-zero*
inf.sup-monoid.add-assoc inf.sup-monoid.add-commute)
also have $\dots = p[?px \sqcap -?ptpy \sqcap y \mapsto ?p2t]$
using 1 *path-splitting-invariant-def* **by** *auto*
also have $\dots = ?pyppy$
using 4 by *auto*
finally show *?thesis*
 \cdot
qed
show *univalent p0 total p0 acyclic (p0 - 1)*
using 1 *path-splitting-invariant-def* **by** *auto*
qed
let $?s = \{ z . \text{regular } z \wedge z \leq p^{T^*} * y \}$
let $?t = \{ z . \text{regular } z \wedge z \leq ?pyppy^{T^*} * (p[[y]]) \}$
have $?pyppy^{T^*} * (p[[y]]) \leq p^{+T^*} * (p[[y]])$
using 1 *path-splitting-invariant-def update-square-plus conv-order*
mult-left-isotone star-isotone **by** *force*
also have $\dots = p^{T^*} * p^T * y$
by (*simp add: conv-plus-commute star.left-plus-circ mult-assoc*)
also have $\dots = p^{T^+} * y$
by (*simp add: star-plus*)
finally have 11: $?pyppy^{T^*} * (p[[y]]) \leq p^{T^+} * y$
 \cdot
hence $?pyppy^{T^*} * (p[[y]]) \leq p^{T^*} * y$
using *mult-left-isotone order-lesseq-imp star.left-plus-below-circ* **by** *blast*
hence 12: $?t \subseteq ?s$
using *order-trans* **by** *auto*
have 13: $y \in ?s$

```

    using 1 bijective-regular path-compression-1b path-splitting-invariant-def by
force
    have 14:  $\neg y \in ?t$ 
    proof
      assume  $y \in ?t$ 
      hence  $y \leq ?pypppy^{T^*} * (p[[y]])$ 
      by simp
      hence  $y \leq p^{T^+} * y$ 
      using 11 dual-order.trans by blast
      hence  $y = \text{root } p \ y$ 
      using 1 find-set-precondition-def loop-root-2 path-splitting-invariant-def by
blast
      hence  $y = ?py$ 
      using 1 by (metis find-set-precondition-def root-successor-loop
path-splitting-invariant-def)
      thus False
      using 1 by simp
    qed
  show  $\text{card } ?t < \text{card } ?s$ 
  apply (rule psubset-card-mono)
  subgoal using finite-regular by simp
  subgoal using 12 13 14 by auto
  done
qed
qed

```

lemma *path-splitting-3*:

path-splitting-invariant $p \ x \ y \ p0 \wedge y = p[[y]] \implies \text{path-splitting-postcondition } p \ x \ y \ p0$

proof –

```

    assume 1: path-splitting-invariant  $p \ x \ y \ p0 \wedge y = p[[y]]$ 
    show path-splitting-postcondition  $p \ x \ y \ p0$ 
    proof (unfold path-splitting-postcondition-def
path-compression-precondition-def, intro conjI)
      show univalent  $p \ \text{total } p \ \text{acyclic } (p - 1)$ 
      using 1 find-set-precondition-def path-splitting-invariant-def by blast+
      show vector  $x \ \text{injective } x \ \text{surjective } x$ 
      using 1 find-set-precondition-def path-splitting-invariant-def by blast+
      show 2: vector  $y \ \text{injective } y \ \text{surjective } y$ 
      using 1 path-splitting-invariant-def by blast+
      show 3:  $p \sqcap 1 = p0 \sqcap 1$ 
      using 1 path-splitting-invariant-aux(4) by blast
      show 4:  $\text{fc } p = \text{fc } p0$ 
      using 1 path-splitting-invariant-aux(5) by blast
      have  $y \leq p0^{T^*} * x$ 
      using 1 path-splitting-invariant-def by simp
      hence 5:  $y * x^T \leq \text{fc } p0$ 
      using 1 by (metis dual-order.trans fc-wcc find-set-precondition-def
shunt-bijective star.circ-decompose-11 star-decompose-1 star-outer-increasing)
    qed
  
```

```

path-splitting-invariant-def)
  have 6:  $y = p0[[y]]$ 
    using 1 path-splitting-invariant-aux(1) by auto
  hence  $y = \text{root } p0 \ y$ 
    using 2 loop-root by auto
  also have  $\dots = \text{root } p0 \ x$ 
    using 1 2 5 find-set-precondition-def path-splitting-invariant-def
same-component-same-root by auto
  also have  $\dots = \text{root } p \ x$ 
    using 1 3 4 by (metis find-set-precondition-def path-splitting-invariant-def
same-root)
  finally show  $y = \text{root } p \ x$ 
  .
  have  $p0^{T^*} * y = y$ 
    using 6 order.antisym path-compression-1b star-left-induct-mult-equal by
auto
  hence 7:  $p0[p0^{T^*} * x - y \longrightarrow (p0 * p0)^T] = p$ 
    using 1 path-splitting-invariant-def by auto
  have  $(p0 * p0)^T * y = y$ 
    using 6 mult-assoc conv-dist-comp by auto
  hence  $y \sqcap p0 * p0 = y \sqcap p0$ 
    using 2 6 by (metis update-postcondition)
  hence 8:  $y \sqcap p = y \sqcap p0 * p0$ 
    using 1 2 6 by (smt update-postcondition)
  have  $p0[p0^{T^*} * x \longrightarrow (p0 * p0)^T] = (p0[p0^{T^*} * x - y \longrightarrow (p0 * p0)^T])[p0^{T^*} *
x \sqcap y \longrightarrow (p0 * p0)^T]$ 
    using 1 bijective-regular path-splitting-invariant-def update-split by blast
  also have  $\dots = p[p0^{T^*} * x \sqcap y \longrightarrow (p0 * p0)^T]$ 
    using 7 by simp
  also have  $\dots = p$ 
    apply (rule update-same-sub)
    using 8 apply simp
    apply simp
    using 1 bijective-regular inf.absorb2 path-splitting-invariant-def by auto
  finally show  $p0[p0^{T^*} * x \longrightarrow (p0 * p0)^T] = p$ 
  .
qed
qed

```

theorem *find-path-splitting*:

```

VARS  $p \ t \ y$ 
[ find-set-precondition  $p \ x \wedge p0 = p$  ]
 $y := x$ ;
WHILE  $y \neq p[[y]]$ 
  INV { path-splitting-invariant  $p \ x \ y \ p0$  }
  VAR {  $(p^{T^*} * y) \downarrow$  }
  DO  $t := p[[y]]$ ;
     $p[y] := p[[p[[y]]]]$ ;
     $y := t$ 

```



```

    OD
  [ path-splitting-postcondition p x y p0 ]
  apply vcg-tc-simp
    apply (fact path-splitting-1)
    apply (fact path-splitting-2)
  by (fact path-splitting-3)

end

```

7 Verifying Union by Rank

In this section we verify the union-by-rank operation of disjoint-set forests. The rank of a node is an upper bound of the height of the subtree rooted at that node. The rank array of a disjoint-set forest maps each node to its rank. This can be represented as a homogeneous relation since the possible rank values are $0, \dots, n-1$ where n is the number of nodes of the disjoint-set forest.

7.1 Peano structures

Since ranks are natural numbers we start by introducing basic Peano arithmetic. Numbers are represented as (relational) points. Constant Z represents the number 0. Constant S represents the successor function. The successor of a number x is obtained by the relational composition $S^T * x$. The composition $S * x$ results in the predecessor of x .

```

class peano-signature =
  fixes Z :: 'a
  fixes S :: 'a

```

The numbers will be used in arrays, which are represented by homogeneous finite relations. Such relations can only represent finitely many numbers. This means that we weaken the Peano axioms, which are usually used to obtain (infinitely many) natural numbers. Axiom *Z-point* specifies that 0 is a number. Axiom *S-univalent* specifies that every number has at most one ‘successor’. Together with axiom *S-total*, which is added later, this means that every number has exactly one ‘successor’. Axiom *S-injective* specifies that numbers with the same successor are equal. Axiom *S-star-Z-top* specifies that every number can be obtained from 0 by finitely many applications of the successor. We omit the Peano axiom $S * Z = bot$ which would specify that 0 is not the successor of any number. Since only finitely many numbers will be represented, the remaining axioms will model successor modulo m for some m depending on the carrier of the algebra. That is, the algebra will be able to represent numbers $0, \dots, m-1$ where the successor of $m-1$ is 0.

```

class skra-peano-1 = stone-kleene-relation-algebra +
stone-relation-algebra-tarski-consistent + peano-signature +
  assumes Z-point: point Z
  assumes S-univalent: univalent S
  assumes S-injective: injective S
  assumes S-star-Z-top:  $S^{T^*} * Z = top$ 
begin

```

```

lemma conv-Z-Z:
   $Z^T * Z = top$ 
  by (simp add: Z-point point-conv-comp)

```

Theorem 9.2

```

lemma Z-below-S-star:
   $Z \leq S^*$ 
proof –
  have top *  $Z^T \leq S^{T^*}$ 
    using S-star-Z-top Z-point shunt-bijection by blast
  thus ?thesis
    using Z-point conv-order conv-star-commute vector-conv-covector by force
qed

```

Theorem 9.3

```

lemma S-connected:
   $S^{T^*} * S^* = top$ 
  by (metis Z-below-S-star S-star-Z-top mult-left-dist-sup sup.orderE sup-commute
top.extremum)

```

Theorem 9.4

```

lemma S-star-connex:
   $S^* \sqcup S^{T^*} = top$ 
  using S-connected S-univalent cancel-separate-eq sup-commute by auto

```

Theorem 9.5

```

lemma Z-sup-conv-S-top:
   $Z \sqcup S^T * top = top$ 
  using S-star-Z-top star.circ-loop-fixpoint sup-commute by auto

```

```

lemma top-S-sup-conv-Z:
   $top * S \sqcup Z^T = top$ 
  by (metis S-star-Z-top conv-dist-comp conv-involutive conv-star-commute
star.circ-back-loop-fixpoint symmetric-top-closed)

```

Theorem 9.1

```

lemma S-inf-1-below-Z:
   $S \sqcap 1 \leq Z$ 
proof –
  have  $(S \sqcap 1) * S^T \leq S \sqcap 1$ 
    by (metis S-injective conv-dist-comp coreflexive-symmetric inf.boundedI
inf.cobounded1 inf.cobounded2 injective-codomain)

```

hence $(S \sqcap 1) * S^{T*} \leq S \sqcap 1$
using *star-right-induct-mult* **by** *blast*
hence $(S \sqcap 1) * S^{T*} * Z \leq (S \sqcap 1) * Z$
by (*simp add: mult-left-isotone*)
also have $\dots \leq Z$
by (*metis comp-left-subdist-inf inf.boundedE mult-1-left*)
finally show *?thesis*
using *S-star-Z-top inf.order-trans top-right-mult-increasing mult-assoc* **by**
auto
qed

lemma *S-inf-1-below-conv-Z*:
 $S \sqcap 1 \leq Z^T$
using *S-inf-1-below-Z conv-order coreflexive-symmetric* **by** *fastforce*

The successor operation provides a convenient way to compare two natural numbers. Namely, $k < m$ if m can be reached from k by finitely many applications of the successor, formally $m \leq S^{T*} * k$ or $k \leq S^* * m$. This does not work for numbers modulo m since comparison depends on the chosen representative. We therefore work with a modified successor relation S' , which is a partial function that computes the successor for all numbers except $m - 1$. If S is surjective, the point M representing the greatest number $m - 1$ is the predecessor of 0 under S . If S is not surjective (like for the set of all natural numbers), $M = \text{bot}$.

abbreviation $S' \equiv S - Z^T$
abbreviation $M \equiv S * Z$

Theorem 11.1

lemma *M-point-iff-S-surjective*:

point $M \longleftrightarrow$ *surjective* S

proof

assume *1: point* M

hence $1 \leq Z^T * S^T * S * Z$

using *comp-associative conv-dist-comp surjective-var* **by** *auto*

hence $Z \leq S^T * S * Z$

using *1 Z-point bijective-reverse mult-assoc* **by** *auto*

also have $\dots \leq S^T * \text{top}$

by (*simp add: comp-isotone mult-assoc*)

finally have $S^T * S^T * \text{top} \sqcup Z \leq S^T * \text{top}$

using *mult-isotone mult-assoc* **by** *force*

hence $S^{T*} * Z \leq S^T * \text{top}$

by (*simp add: star-left-induct mult-assoc*)

thus *surjective* S

by (*simp add: S-star-Z-top order.antisym surjective-conv-total*)

next

assume *surjective* S

thus *point* M

by (*metis S-injective Z-point comp-associative injective-mult-closed*)

qed

[Theorem 10.1](#)

lemma *S'*-univalent:

univalent S'

by (*simp add: S-univalent univalent-inf-closed*)

[Theorem 10.2](#)

lemma *S'*-injective:

injective S'

by (*simp add: S-injective injective-inf-closed*)

[Theorem 10.9](#)

lemma *S'*-Z:

$S' * Z = \text{bot}$

by (*simp add: Z-point covector-vector-comp injective-comp-right-dist-inf*)

[Theorem 10.4](#)

lemma *S'*-irreflexive:

irreflexive S'

using *S-inf-1-below-conv-Z order-lesseq-imp p-shunting-swap pp-increasing* **by**
blast

end

class *skra-peano-2* = *skra-peano-1* +

assumes *S-total: total S*

begin

lemma *S*-mapping:

mapping S

by (*simp add: S-total S-univalent*)

[Theorem 11.2](#)

lemma *M-bot-iff-S-not-surjective*:

$M \neq \text{bot} \longleftrightarrow \text{surjective } S$

proof

assume $M \neq \text{bot}$

hence $\text{top} * S * Z = \text{top}$

by (*metis S-mapping Z-point bijective-regular comp-associative
mapping-regular regular-mult-closed tarski*)

hence $Z^T \leq \text{top} * S$

using *M-point-iff-S-surjective S-injective Z-point comp-associative
injective-mult-closed* **by** *auto*

thus *surjective S*

using *sup.orderE top-S-sup-conv-Z* **by** *fastforce*

next

assume *surjective S*

thus $M \neq \text{bot}$

using *M-point-iff-S-surjective consistent covector-bot-closed* by force
qed

Theorem 11.3

lemma *M-point-or-bot*:

point $M \vee M = \text{bot}$

using *M-bot-iff-S-not-surjective M-point-iff-S-surjective* by blast

Alternative way to express S'

Theorem 12.1

lemma *S'-var*:

$S' = S - M$

proof –

have $S' = S * (1 - Z^T)$

by (*simp add: Z-point covector-comp-inf vector-conv-compl*)

also have $\dots = S * (1 - Z)$

by (*metis conv-complement one-inf-conv*)

also have $\dots = S * 1 \sqcap S * -Z$

by (*simp add: S-mapping univalent-comp-left-dist-inf*)

also have $\dots = S - M$

by (*simp add: comp-mapping-complement S-mapping*)

finally show *?thesis*

qed

Special case of just 1 number

Theorem 12.2

lemma *M-is-Z-iff-1-is-top*:

$M = Z \longleftrightarrow 1 = \text{top}$

proof

assume $M = Z$

hence $Z = S^T * Z$

by (*metis S-mapping Z-point order.antisym bijective-reverse inf.eq-refl shunt-mapping*)

thus $1 = \text{top}$

by (*metis S-star-Z-top Z-point inf.eq-refl star-left-induct sup.absorb2 symmetric-top-closed top-le*)

next

assume $1 = \text{top}$

thus $M = Z$

using *S-mapping comp-right-one mult-1-left* by auto

qed

Theorem 12.3

lemma *S-irreflexive*:

assumes $M \neq Z$

shows *irreflexive S*

proof –

```

have  $(S \sqcap 1) * S^T \leq S \sqcap 1$ 
  by (smt (z3) S-injective S-mapping coreflexive-comp-top-inf dual-order.eq-iff
inf.cobounded1 inf.sup-monoid.add-commute inf.sup-same-context
mult-left-isotone one-inf-conv top-right-mult-increasing total-var)
  hence  $(S \sqcap 1) * S^{T*} \leq S \sqcap 1$ 
    using star-right-induct-mult by blast
  hence  $(S \sqcap 1) * S^{T*} * Z \leq (S \sqcap 1) * Z$ 
    by (simp add: mult-left-isotone)
  also have  $\dots = M \sqcap Z$ 
    by (simp add: Z-point injective-comp-right-dist-inf)
  also have  $\dots = \text{bot}$ 
    by (smt (verit, cfv-threshold) M-point-or-bot assms Z-point
bijjective-one-closed bijjective-regular comp-associative conv-complement
coreflexive-comp-top-inf epm-3 inf.sup-monoid.add-commute one-inf-conv
regular-mult-closed star.circ-increasing star.circ-zero tarski vector-conv-covector
vector-export-comp-unit)
  finally have  $S \sqcap 1 \leq \text{bot}$ 
    using S-star-Z-top comp-associative le-bot top-right-mult-increasing by
fastforce
  thus ?thesis
    using le-bot pseudo-complement by blast
qed

```

We show that S' satisfies most properties of S .

```

lemma M-regular:
  regular M
  using S-mapping Z-point bijjective-regular mapping-regular regular-mult-closed
by blast

```

```

lemma S'-regular:
  regular S'
  using S-mapping mapping-regular by auto

```

Theorem 10.3

```

lemma S'-star-Z-top:
   $S^{T*} * Z = \text{top}$ 
proof –
  have  $S^{T*} * Z = (S' \sqcup (S \sqcap M))^{T*} * Z$ 
    by (metis M-regular maddux-3-11-pp S'-var)
  also have  $\dots \leq S^{T*} * Z$ 
proof (cases M = bot)
  case True
    thus ?thesis
    by simp
next
  case False
  hence point M
    using M-point-or-bot by auto
  hence arc (S \sqcap M)

```

using *S*-mapping mapping-inf-point-arc **by** *blast*
hence $1: \text{arc } ((S \sqcap M)^T)$
using *conv-involutive* **by** *auto*
have $2: S \sqcap M \leq Z^T$
by (*metis S'-var Z-point bijective-regular conv-complement inf.cobounded2*
p-shunting-swap)
have $(S' \sqcup (S \sqcap M))^{T*} * Z = (S'^T \sqcup (S \sqcap M)^T)^* * Z$
by (*simp add: S'-var conv-dist-sup*)
also have $\dots = (S'^{T*} * (S \sqcap M)^T * S'^{T*} \sqcup S'^{T*}) * Z$
using *1 star-arc-decompose sup-commute* **by** *auto*
also have $\dots = S'^{T*} * (S \sqcap M)^T * S'^{T*} * Z \sqcup S'^{T*} * Z$
using *mult-right-dist-sup* **by** *auto*
also have $\dots \leq S'^{T*} * Z^{TT} * S'^{T*} * Z \sqcup S'^{T*} * Z$
using *2* **by** (*meson comp-isotone conv-isotone inf.eq-refl semiring.add-mono*)
also have $\dots \leq S'^{T*} * Z$
by (*metis Z-point comp-associative conv-involutive le-supI mult-right-isotone*
top.extremum)
finally show *?thesis*
qed
finally show *?thesis*
using *S-star-Z-top top-le* **by** *auto*
qed

Theorem 10.5

lemma *Z-below-S'-star*:

$$Z \leq S'^*$$

by (*metis S'-star-Z-top Z-point comp-associative comp-right-one conv-order*
conv-star-commute mult-right-isotone vector-conv-covector)

Theorem 10.6

lemma *S'-connected*:

$$S'^{T*} * S'^* = \text{top}$$

by (*metis Z-below-S'-star S'-star-Z-top mult-left-dist-sup sup.orderE*
sup-commute top.extremum)

Theorem 10.7

lemma *S'-star-connex*:

$$S'^* \sqcup S'^{T*} = \text{top}$$

using *S'-connected S'-univalent cancel-separate-eq sup-commute* **by** *auto*

Theorem 10.8

lemma *Z-sup-conv-S'-top*:

$$Z \sqcup S'^T * \text{top} = \text{top}$$

using *S'-star-Z-top star.circ-loop-fixpoint sup-commute* **by** *auto*

lemma *top-S'-sup-conv-Z*:

$$\text{top} * S' \sqcup Z^T = \text{top}$$

by (*metis S'-star-Z-top conv-dist-comp conv-involutive conv-star-commute*
star.circ-back-loop-fixpoint symmetric-top-closed)

end

7.2 Initialising Ranks

We show that the rank array satisfies three properties which are established/preserved by the union-find operations. First, every node has a rank, that is, the rank array is a mapping. Second, the rank of a node is strictly smaller than the rank of its parent, except if the node is a root. This implies that the rank of a node is an upper bound on the height of its subtree. Third, the number of roots in the disjoint-set forest (the number of disjoint sets) is not larger than $m - k$ where m is the total number of nodes and k is the maximum rank of any node. The third property is useful to show that ranks never overflow (exceed $m - 1$). To compare the number of roots and $m - k$ we use the existence of an injective univalent relation between the set of roots and the set of $m - k$ largest numbers, both represented as vectors. The three properties are captured in *rank-property*.

```
class skra-peano-3 = stone-kleene-relation-algebra-tarski-finite-regular +
skra-peano-2
begin
```

```
definition card-less-eq v w ≡ ∃ i . injective i ∧ univalent i ∧ regular i ∧ v ≤ i * w
definition rank-property p rank ≡ mapping rank ∧ (p - 1) * rank ≤ rank * S+
∧ card-less-eq (roots p) (-(S+ * rankT * top))
```

end

```
class skra-peano-4 = stone-kleene-relation-algebra-choose-point-finite-regular +
skra-peano-2
begin
```

```
subclass skra-peano-3 ..
```

The initialisation loop is augmented by setting the rank of each node to 0. The resulting rank array satisfies the desired properties explained above.

theorem *init-ranks*:

```
  VARS h p x rank
  [ True ]
  FOREACH x
  USING h
  INV { p - h = 1 - h ∧ rank - h = ZT - h }
  DO p := make-set p x;
    rank[x] := Z
  OD
  [ p = 1 ∧ disjoint-set-forest p ∧ rank = ZT ∧ rank-property p rank ∧ h = bot ]
proof vcp-tc-simp
  fix h p rank
```



```

let ?x = choose-point h
let ?m = make-set p ?x
let ?rank = rank[?x→Z]
assume 1: regular h ∧ vector h ∧ p - h = 1 - h ∧ rank - h = ZT - h ∧ h ≠
bot
show vector (-?x ⊓ h) ∧
  ?m ⊓ (--?x ⊔ -h) = 1 ⊓ (--?x ⊔ -h) ∧
  ?rank ⊓ (--?x ⊔ -h) = ZT ⊓ (--?x ⊔ -h) ∧
  card { x . regular x ∧ x ≤ -?x ∧ x ≤ h } < h↓
proof (intro conjI)
  show vector (-?x ⊓ h)
    using 1 choose-point-point vector-complement-closed vector-inf-closed by
blast
  have 2: point ?x ∧ regular ?x
    using 1 bijective-regular choose-point-point by blast
  have 3: -h ≤ -?x
    using choose-point-decreasing p-antitone-iff by auto
  have 4: ?x ⊓ ?m = ?x * ?xT ∧ -?x ⊓ ?m = -?x ⊓ p
    using 1 choose-point-point make-set-function make-set-postcondition-def by
auto
  have ?m ⊓ (--?x ⊔ -h) = (?m ⊓ ?x) ⊔ (?m - h)
    using 2 comp-inf.comp-left-dist-sup by auto
  also have ... = ?x * ?xT ⊔ (?m ⊓ -?x ⊓ -h)
    using 3 4 by (smt (z3) inf-absorb2 inf-assoc inf-commute)
  also have ... = ?x * ?xT ⊔ (1 - h)
    using 1 3 4 inf.absorb2 inf.sup-monoid.add-assoc inf-commute by auto
  also have ... = (1 ⊓ ?x) ⊔ (1 - h)
    using 2 by (metis inf.cobounded2 inf.sup-same-context one-inf-conv
vector-covector)
  also have ... = 1 ⊓ (--?x ⊔ -h)
    using 2 comp-inf.semiring.distrib-left by auto
  finally show ?m ⊓ (--?x ⊔ -h) = 1 ⊓ (--?x ⊔ -h)
  .
  have 5: ?x ⊓ ?rank = ?x ⊓ ZT ∧ -?x ⊓ ?rank = -?x ⊓ rank
    by (smt (z3) inf-commute order-refl update-inf-different update-inf-same)
  have ?rank ⊓ (--?x ⊔ -h) = (?rank ⊓ ?x) ⊔ (?rank - h)
    using 2 comp-inf.comp-left-dist-sup by auto
  also have ... = (?x ⊓ ZT) ⊔ (?rank ⊓ -?x ⊓ -h)
    using 3 5 by (smt (z3) inf-absorb2 inf-assoc inf-commute)
  also have ... = (ZT ⊓ ?x) ⊔ (ZT - h)
    using 1 3 5 inf.absorb2 inf.sup-monoid.add-assoc inf-commute by auto
  also have ... = ZT ⊓ (--?x ⊔ -h)
    using 2 comp-inf.semiring.distrib-left by auto
  finally show ?rank ⊓ (--?x ⊔ -h) = ZT ⊓ (--?x ⊔ -h)
  .
  have 5: ¬ ?x ≤ -?x
    using 1 2 by (metis comp-commute-below-diversity conv-order
inf.cobounded2 inf-absorb2 pseudo-complement strict-order-var top.extremum)
  have 6: ?x ≤ h

```

```

    using 1 by (metis choose-point-decreasing)
  show card { x . regular x ∧ x ≤ -?x ∧ x ≤ h } < h↓
  apply (rule psubset-card-mono)
  using finite-regular apply simp
  using 2 5 6 by auto
qed
next
show rank-property 1 (ZT)
proof (unfold rank-property-def, intro conjI)
  show univalent (ZT) total (ZT)
  using Z-point surjective-conv-total by auto
  show (1 - 1) * (ZT) ≤ (ZT) * S+
  by simp
  have top ≤ 1 * -(S+ * Z * top)
  by (simp add: S+-Z comp-associative star-simulation-right-equal)
  thus card-less-eq (roots 1) (-(S+ * ZT * top))
  by (metis conv-involutive inf.idem mapping-one-closed regular-one-closed
card-less-eq-def bijective-one-closed)
qed
qed
end

```

7.3 Union by Rank

We show that path compression and union-by-rank preserve the rank property.

context *stone-kleene-relation-algebra-tarski-finite-regular*
begin

lemma *union-sets-1-swap*:

```

  assumes union-sets-precondition p0 x y
    and path-compression-postcondition p1 x r p0
    and path-compression-postcondition p2 y s p1
  shows union-sets-postcondition (p2[s→r]) x y p0
proof (unfold union-sets-postcondition-def union-sets-precondition-def, intro
conjI)
  let ?p = p2[s→r]
  have 1: disjoint-set-forest p1 ∧ point r ∧ r = root p1 x ∧ p1 ⊔ 1 = p0 ⊔ 1 ∧
fc p1 = fc p0
  using assms(2) path-compression-precondition-def
path-compression-postcondition-def by auto
  have 2: disjoint-set-forest p2 ∧ point s ∧ s = root p2 y ∧ p2 ⊔ 1 = p1 ⊔ 1 ∧
fc p2 = fc p1
  using assms(3) path-compression-precondition-def
path-compression-postcondition-def by auto
  hence 3: fc p2 = fc p0
  using 1 by simp
  show 4: univalent ?p

```

```

    using 1 2 update-univalent by blast
show total ?p
    using 1 2 bijective-regular update-total by blast
show acyclic (?p - 1)
proof (cases r = s)
  case True
  thus ?thesis
    using 2 update-acyclic-5 by fastforce
next
  case False
  hence bot = s  $\sqcap$  r
    using 1 2 distinct-points inf-commute by blast
  also have ... = s  $\sqcap$  p1T* * r
    using 1 by (smt root-transitive-successor-loop)
  also have ... = s  $\sqcap$  p2T* * r
    using 1 2 by (smt (z3) inf-assoc inf-commute same-root)
  finally have r  $\sqcap$  p2* * s = bot
    using schroeder-1 conv-star-commute inf.sup-monoid.add-commute by
fastforce
  thus ?thesis
    using 1 2 update-acyclic-4 by blast
qed
show vector x
  using assms(1) by (simp add: union-sets-precondition-def)
show injective x
  using assms(1) by (simp add: union-sets-precondition-def)
show surjective x
  using assms(1) by (simp add: union-sets-precondition-def)
show vector y
  using assms(1) by (simp add: union-sets-precondition-def)
show injective y
  using assms(1) by (simp add: union-sets-precondition-def)
show surjective y
  using assms(1) by (simp add: union-sets-precondition-def)
show fc ?p = wcc (p0  $\sqcup$  x * yT)
proof (rule order.antisym)
  have s = p2[[s]]
    using 2 by (metis root-successor-loop)
  hence s * sT  $\leq$  p2T
    using 2 eq-refl shunt-bijective by blast
  hence s * sT  $\leq$  p2
    using 2 conv-order coreflexive-symmetric by fastforce
  hence s  $\leq$  p2 * s
    using 2 shunt-bijective by blast
  hence 5: p2[[s]]  $\leq$  s
    using 2 shunt-mapping by blast
  have s  $\sqcap$  p2  $\leq$  s * (top  $\sqcap$  sT * p2)
    using 2 by (metis dedekind-1)
  also have ... = s * sT * p2

```

by (*simp add: mult-assoc*)
 also have $\dots \leq s * s^T$
 using 5 by (*metis comp-associative conv-dist-comp conv-involutive conv-order mult-right-isotone*)
 also have $\dots \leq 1$
 using 2 by *blast*
 finally have 6: $s \sqcap p2 \leq 1$
 by *simp*
 have $p0 \leq wcc\ p0$
 by (*simp add: star.circ-sub-dist-1*)
 also have $\dots = wcc\ p2$
 using 3 by (*simp add: star-decompose-1*)
 also have 7: $\dots \leq wcc\ ?p$
proof –
 have $wcc\ p2 = wcc\ ((-s \sqcap p2) \sqcup (s \sqcap p2))$
 using 2 by (*metis bijective-regular inf.sup-monoid.add-commute maddux-3-11-pp*)
 also have $\dots \leq wcc\ ((-s \sqcap p2) \sqcup 1)$
 using 6 *wcc-isotone sup-right-isotone* by *simp*
 also have $\dots = wcc\ (-s \sqcap p2)$
 using *wcc-with-loops* by *simp*
 also have $\dots \leq wcc\ ?p$
 using *wcc-isotone sup-ge2* by *blast*
finally show *?thesis*
 by *simp*
qed
 finally have 8: $p0 \leq wcc\ ?p$
 by *force*
 have $s \leq p2^{T*} * y$
 using 2 by (*metis inf-le1*)
 hence 9: $s * y^T \leq p2^{T*}$
 using *assms(1) shunt-bijective union-sets-precondition-def* by *blast*
 hence $y * s^T \leq p2^{T*}$
 using *conv-dist-comp conv-order conv-star-commute* by *force*
 also have $\dots \leq wcc\ p2$
 by (*simp add: star.circ-sub-dist*)
 also have $\dots \leq wcc\ ?p$
 using 7 by *simp*
 finally have 10: $y * s^T \leq wcc\ ?p$
 by *simp*
 have 11: $s * r^T \leq wcc\ ?p$
 using 1 2 *star.circ-sub-dist-1 sup-assoc vector-covector* by *auto*
 have $r \leq p1^{T*} * x$
 using 1 by (*metis inf-le1*)
 hence 12: $r * x^T \leq p1^{T*}$
 using *assms(1) shunt-bijective union-sets-precondition-def* by *blast*
 also have $\dots \leq wcc\ p1$
 using *star-isotone sup-ge2* by *blast*
 also have $\dots = wcc\ p2$

using 2 by (*simp add: star-decompose-1*)
 also have $\dots \leq wcc \ ?p$
 using 7 by *simp*
 finally have 13: $r * x^T \leq wcc \ ?p$
 by *simp*
 have $x \leq x * r^T * r \wedge y \leq y * s^T * s$
 using 1 2 *shunt-bijective* by *blast*
 hence $y * x^T \leq y * s^T * s * (x * r^T * r)^T$
 using *comp-isotone conv-isotone* by *blast*
 also have $\dots = y * s^T * s * r^T * r * x^T$
 by (*simp add: comp-associative conv-dist-comp*)
 also have $\dots \leq wcc \ ?p * (s * r^T) * (r * x^T)$
 using 10 by (*metis mult-left-isotone mult-assoc*)
 also have $\dots \leq wcc \ ?p * wcc \ ?p * (r * x^T)$
 using 11 by (*metis mult-left-isotone mult-right-isotone*)
 also have $\dots \leq wcc \ ?p * wcc \ ?p * wcc \ ?p$
 using 13 by (*metis mult-right-isotone*)
 also have $\dots = wcc \ ?p$
 by (*simp add: star.circ-transitive-equal*)
 finally have $x * y^T \leq wcc \ ?p$
 by (*metis conv-dist-comp conv-involutive conv-order wcc-equivalence*)
 hence $p0 \sqcup x * y^T \leq wcc \ ?p$
 using 8 by *simp*
 hence $wcc (p0 \sqcup x * y^T) \leq wcc \ ?p$
 using *wcc-below-wcc* by *simp*
 thus $wcc (p0 \sqcup x * y^T) \leq fc \ ?p$
 using 4 *fc-wcc* by *simp*
 have $-s \sqcap p2 \leq wcc \ p2$
 by (*simp add: inf.coboundedI2 star.circ-sub-dist-1*)
 also have $\dots = wcc \ p0$
 using 3 by (*simp add: star-decompose-1*)
 also have $\dots \leq wcc (p0 \sqcup y * x^T)$
 by (*simp add: wcc-isotone*)
 finally have 14: $-s \sqcap p2 \leq wcc (p0 \sqcup y * x^T)$
 by *simp*
 have $s * y^T \leq wcc \ p2$
 using 9 *inf.order-trans star.circ-sub-dist sup-commute* by *fastforce*
 also have $\dots = wcc \ p0$
 using 1 2 by (*simp add: star-decompose-1*)
 also have $\dots \leq wcc (p0 \sqcup y * x^T)$
 by (*simp add: wcc-isotone*)
 finally have 15: $s * y^T \leq wcc (p0 \sqcup y * x^T)$
 by *simp*
 have 16: $y * x^T \leq wcc (p0 \sqcup y * x^T)$
 using *le-supE star.circ-sub-dist-1* by *blast*
 have $x * r^T \leq p1^*$
 using 12 *conv-dist-comp conv-order conv-star-commute* by *fastforce*
 also have $\dots \leq wcc \ p1$
 using *star.circ-sub-dist sup-commute* by *fastforce*

also have $\dots = \text{wcc } p0$
using *1* **by** (*simp add: star-decompose-1*)
also have $\dots \leq \text{wcc } (p0 \sqcup y * x^T)$
by (*simp add: wcc-isotone*)
finally have *17*: $x * r^T \leq \text{wcc } (p0 \sqcup y * x^T)$
by *simp*
have $r \leq r * x^T * x \wedge s \leq s * y^T * y$
using *assms(1)* *shunt-bijective union-sets-precondition-def* **by** *blast*
hence $s * r^T \leq s * y^T * y * (r * x^T * x)^T$
using *comp-isotone conv-isotone* **by** *blast*
also have $\dots = s * y^T * y * x^T * x * r^T$
by (*simp add: comp-associative conv-dist-comp*)
also have $\dots \leq \text{wcc } (p0 \sqcup y * x^T) * (y * x^T) * (x * r^T)$
using *15* **by** (*metis mult-left-isotone mult-assoc*)
also have $\dots \leq \text{wcc } (p0 \sqcup y * x^T) * \text{wcc } (p0 \sqcup y * x^T) * (x * r^T)$
using *16* **by** (*metis mult-left-isotone mult-right-isotone*)
also have $\dots \leq \text{wcc } (p0 \sqcup y * x^T) * \text{wcc } (p0 \sqcup y * x^T) * \text{wcc } (p0 \sqcup y * x^T)$
using *17* **by** (*metis mult-right-isotone*)
also have $\dots = \text{wcc } (p0 \sqcup y * x^T)$
by (*simp add: star.circ-transitive-equal*)
finally have $?p \leq \text{wcc } (p0 \sqcup y * x^T)$
using *1 2 14* *vector-covector* **by** *auto*
hence $\text{wcc } ?p \leq \text{wcc } (p0 \sqcup y * x^T)$
using *wcc-below-wcc* **by** *blast*
also have $\dots = \text{wcc } (p0 \sqcup x * y^T)$
using *conv-dist-comp conv-dist-sup sup-assoc sup-commute* **by** *auto*
finally show $fc ?p \leq \text{wcc } (p0 \sqcup x * y^T)$
using *4* *fc-wcc* **by** *simp*

qed
qed

lemma *union-sets-1-skip*:

assumes *union-sets-precondition* $p0$ x y
and *path-compression-postcondition* $p1$ x r $p0$
and *path-compression-postcondition* $p2$ y r $p1$
shows *union-sets-postcondition* $p2$ x y $p0$
proof (*unfold union-sets-postcondition-def union-sets-precondition-def, intro conjI*)
have *1*: $\text{point } r \wedge r = \text{root } p1 \ x \wedge fc \ p1 = fc \ p0 \wedge \text{disjoint-set-forest } p2 \wedge r = \text{root } p2 \ y \wedge fc \ p2 = fc \ p1$
using *assms(2,3)* *path-compression-precondition-def*
path-compression-postcondition-def **by** *auto*
thus *univalent* $p2$ *total* $p2$ *acyclic* $(p2 - 1)$
by *auto*
show *vector* x *injective* x *surjective* x *vector* y *injective* y *surjective* y
using *assms(1)* *union-sets-precondition-def* **by** *auto*
have $r \leq p1^{T*} * x$
using *1* **by** (*metis inf-le1*)
hence $r * x^T \leq p1^{T*}$

```

    using assms(1) shunt-bijective union-sets-precondition-def by blast
  hence 2:  $x * r^T \leq p1^*$ 
    using conv-dist-comp conv-order conv-star-commute by force
  have  $r \leq p2^{T*} * y$ 
    using 1 by (metis inf-le1)
  hence 3:  $r * y^T \leq p2^{T*}$ 
    using assms(1) shunt-bijective union-sets-precondition-def by blast
  have  $x * y^T \leq x * r^T * r * y^T$ 
    using 1 mult-left-isotone shunt-bijective by blast
  also have  $\dots \leq p1^* * p2^{T*}$ 
    using 2 3 by (metis comp-associative comp-isotone)
  also have  $\dots \leq wcc\ p0$ 
    using 1 by (metis star.circ-mult-upper-bound star-decompose-1 star-isotone
sup-ge2 star.circ-sub-dist)
  finally show  $fc\ p2 = wcc\ (p0 \sqcup x * y^T)$ 
    using 1 by (smt (z3) fc-star star-decompose-1 sup-absorb1 wcc-sup-wcc
star.circ-sub-dist-3 sup-commute wcc-equivalence)
qed

end

```

syntax

```

-Cond1 :: 'bexp  $\Rightarrow$  'com  $\Rightarrow$  'com ((IIF -/ THEN -/ FI) [0,0] 61)
translations IF b THEN c FI == IF b THEN c ELSE SKIP FI

```

context *skra-peano-3*

begin

lemma *path-compression-preserves-rank-property*:

```

  assumes path-compression-postcondition p x y p0
    and disjoint-set-forest p0
    and rank-property p0 rank
  shows rank-property p rank
proof (unfold rank-property-def, intro conjI)
  let  $?px = p0^{T*} * x$ 
  have 1: point y
    using assms(1,2) path-compression-postcondition-def
path-compression-precondition-def root-point by auto
  have 2: vector  $?px$ 
    using assms(1) comp-associative path-compression-postcondition-def
path-compression-precondition-def by auto
  have root p0 x = root p x
    by (smt (verit) assms(1,2) path-compression-postcondition-def
path-compression-precondition-def same-root)
  hence root p0 x = y
    using assms(1) path-compression-postcondition-def
path-compression-precondition-def by auto
  hence  $?px \leq p0^* * y$ 
    by (meson assms(1,2) path-splitting-invariant-aux-1(3))

```

path-compression-precondition-def path-compression-postcondition-def
hence $?px * y^T \leq p0^*$
using *1 shunt-bijective* **by** *blast*
hence $?px \sqcap y^T \leq p0^*$
using *1 2* **by** (*simp add: vector-covector*)
also have $\dots = (p0 - 1)^+ \sqcup 1$
using *reachable-without-loops star-left-unfold-equal sup-commute* **by** *fastforce*
finally have $?3: ?px \sqcap y^T \sqcap -1 \leq (p0 - 1)^+$
using *half-shunting* **by** *blast*
have $p0[?px \rightarrow y] = p$
using *assms(1) path-compression-postcondition-def* **by** *simp*
hence $(p - 1) * rank = (?px \sqcap y^T \sqcap -1) * rank \sqcup (-?px \sqcap p0 \sqcap -1) * rank$
using *inf-sup-distrib2 mult-right-dist-sup* **by** *force*
also have $\dots \leq (?px \sqcap y^T \sqcap -1) * rank \sqcup (p0 - 1) * rank$
by (*meson comp-inf.mult-semi-associative le-infE mult-left-isotone*
sup-right-isotone)
also have $\dots \leq (?px \sqcap y^T \sqcap -1) * rank \sqcup rank * S'^+$
using *assms(3) rank-property-def sup-right-isotone* **by** *auto*
also have $\dots \leq (p0 - 1)^+ * rank \sqcup rank * S'^+$
using *3 mult-left-isotone sup-left-isotone* **by** *blast*
also have $\dots \leq rank * S'^+$
proof –
have $(p0 - 1)^* * rank \leq rank * S'^*$
using *assms(3) rank-property-def star-simulation-left star.left-plus-circ* **by**
fastforce
hence $(p0 - 1)^+ * rank \leq (p0 - 1) * rank * S'^*$
by (*simp add: comp-associative mult-right-isotone*)
also have $\dots \leq rank * S'^+$
by (*smt (z3) assms(3) rank-property-def comp-associative*
comp-left-subdist-inf inf.boundedE inf.sup-right-divisibility
star.circ-transitive-equal)
finally show *?thesis*
by *simp*
qed
finally show $(p - 1) * rank \leq rank * S'^+$
show *univalent rank total rank*
using *rank-property-def assms(3)* **by** *auto*
show *card-less-eq (roots p) (-(S'^+ * rank^T * top))*
using *assms(1,3) path-compression-postcondition-def rank-property-def* **by**
auto
qed

theorem *union-sets-by-rank:*

VARs p r s rank

[*union-sets-precondition p x y* \wedge *rank-property p rank* \wedge $p0 = p$]

$r :=$ *find-set p x*;

$p :=$ *path-compression p x r*;

$s :=$ *find-set p y*;


```

p := path-compression p y s;
IF r ≠ s THEN
  IF rank[[r]] ≤  $S'^+$  * (rank[[s]]) THEN
    p[[r]] := s
  ELSE
    p[[s]] := r;
    IF rank[[r]] = rank[[s]] THEN
      rank[[r]] :=  $S'^T$  * (rank[[r]])
    FI
  FI
FI
[ union-sets-postcondition p x y p0 ∧ rank-property p rank ]
proof vcg-tc-simp
  fix rank
  let ?r = find-set p0 x
  let ?p1 = path-compression p0 x ?r
  let ?s = find-set ?p1 y
  let ?p2 = path-compression ?p1 y ?s
  let ?p5 = path-compression ?p1 y ?r
  let ?rr = rank[[?r]]
  let ?rs = rank[[?s]]
  let ?rank = rank[[?r] →  $S'^T$  * ?rs]
  let ?p3 = ?p2[[?r] → ?s]
  let ?p4 = ?p2[[?s] → ?r]
  assume 1: union-sets-precondition p0 x y ∧ rank-property p0 rank
  hence 2: path-compression-postcondition ?p1 x ?r p0
    using find-set-function find-set-postcondition-def find-set-precondition-def
path-compression-function path-compression-precondition-def
union-sets-precondition-def by auto
  hence 3: path-compression-postcondition ?p2 y ?s ?p1
    using 1 find-set-function find-set-postcondition-def find-set-precondition-def
path-compression-function path-compression-precondition-def
union-sets-precondition-def path-compression-postcondition-def by meson
  have rank-property ?p1 rank
    using 1 2 path-compression-preserves-rank-property
union-sets-precondition-def by blast
  hence 4: rank-property ?p2 rank
    using 1 2 3 by (meson path-compression-preserves-rank-property
path-compression-postcondition-def path-compression-precondition-def)
  have 5: point ?r point ?s
    using 2 3 path-compression-postcondition-def
path-compression-precondition-def by auto
  hence 6: point ?rr point ?rs
    using 1 comp-associative read-injective read-surjective rank-property-def by
auto
  have top ≤  $S'^*$  ⊔  $S'^+T$ 
    by (metis  $S'$ -star-connex conv-dist-comp conv-star-commute eq-refl
star.circ-reflexive star-left-unfold-equal star-simulation-right-equal sup.orderE
sup-monoid.add-assoc)

```

hence 7: $-S'^{+T} \leq S'^*$
by (*metis comp-inf.case-split-left comp-inf.star.circ-plus-one*
comp-inf.star.circ-sup-2 half-shunting)
show $(?r \neq ?s \longrightarrow (?rr \leq S'^+ * ?rs \longrightarrow \text{union-sets-postcondition } ?p3 \ x \ y \ p0 \ \wedge$
rank-property } ?p3 \ rank) \ \wedge
 $(\neg ?rr \leq S'^+ * ?rs \longrightarrow ((?rr = ?rs \longrightarrow$
union-sets-postcondition } ?p4 \ x \ y \ p0 \ \wedge \ \text{rank-property } ?p4 \ ?rank) \ \wedge
 $(?rr \neq ?rs \longrightarrow \text{union-sets-postcondition } ?p4 \ x$
y \ p0 \ \wedge \ \text{rank-property } ?p4 \ rank)))) \ \wedge
 $(?r = ?s \longrightarrow \text{union-sets-postcondition } ?p5 \ x \ y \ p0 \ \wedge \ \text{rank-property } ?p5 \ rank)$
proof
show $?r \neq ?s \longrightarrow (?rr \leq S'^+ * ?rs \longrightarrow \text{union-sets-postcondition } ?p3 \ x \ y \ p0$
 $\wedge \ \text{rank-property } ?p3 \ rank) \ \wedge$
 $(\neg ?rr \leq S'^+ * ?rs \longrightarrow ((?rr = ?rs \longrightarrow$
union-sets-postcondition } ?p4 \ x \ y \ p0 \ \wedge \ \text{rank-property } ?p4 \ ?rank) \ \wedge
 $(?rr \neq ?rs \longrightarrow \text{union-sets-postcondition } ?p4 \ x$
y \ p0 \ \wedge \ \text{rank-property } ?p4 \ rank))))
proof
assume 8: $?r \neq ?s$
show $(?rr \leq S'^+ * ?rs \longrightarrow \text{union-sets-postcondition } ?p3 \ x \ y \ p0 \ \wedge$
rank-property } ?p3 \ rank) \ \wedge
 $(\neg ?rr \leq S'^+ * ?rs \longrightarrow ((?rr = ?rs \longrightarrow \text{union-sets-postcondition } ?p4 \ x$
y \ p0 \ \wedge \ \text{rank-property } ?p4 \ ?rank) \ \wedge
 $(?rr \neq ?rs \longrightarrow \text{union-sets-postcondition } ?p4 \ x \ y \ p0 \ \wedge$
rank-property } ?p4 \ rank))))
proof
show $?rr \leq S'^+ * ?rs \longrightarrow \text{union-sets-postcondition } ?p3 \ x \ y \ p0 \ \wedge$
rank-property } ?p3 \ rank
proof
assume 9: $?rr \leq S'^+ * ?rs$
show *union-sets-postcondition } ?p3 \ x \ y \ p0 \ \wedge \ \text{rank-property } ?p3 \ rank*
proof
show *union-sets-postcondition } ?p3 \ x \ y \ p0*
using 1 2 3 by (*simp add: union-sets-1*)
show *rank-property } ?p3 \ rank*
proof (*unfold rank-property-def, intro conjI*)
show *univalent rank total rank*
using 1 rank-property-def by auto
have $?s \leq -?r$
using 5 8 by (*meson order.antisym bijective-regular*
point-in-vector-or-complement point-in-vector-or-complement-2)
hence $?r \sqcap ?s^T \sqcap 1 = \text{bot}$
by (*metis (full-types) bot-least inf.left-commute*
inf.sup-monoid.add-commute one-inf-conv pseudo-complement)
hence $?p3 \sqcap 1 \leq ?p2$
by (*smt half-shunting inf.cobounded2 pseudo-complement*
regular-one-closed semiring.add-mono sup-commute)
hence $\text{roots } ?p3 \leq \text{roots } ?p2$
by (*simp add: mult-left-isotone*)

```

thus card-less-eq (roots ?p3) (-(S'+ * rankT * top))
  using 4 by (meson rank-property-def card-less-eq-def order-trans)
have (?p3 - 1) * rank = (?r □ ?sT □ -1) * rank □ (-?r □ ?p2 □
-1) * rank
  using comp-inf.semiring.distrib-right mult-right-dist-sup by auto
also have ... ≤ (?r □ ?sT □ -1) * rank □ (?p2 - 1) * rank
  using comp-inf.mult-semi-associative mult-left-isotone
sup-right-isotone by auto
also have ... ≤ (?r □ ?sT □ -1) * rank □ rank * S'+
  using 4 sup-right-isotone rank-property-def by blast
also have ... ≤ (?r □ ?sT) * rank □ rank * S'+
  using inf-le1 mult-left-isotone sup-left-isotone by blast
also have ... = ?r * ?sT * rank □ rank * S'+
  using 5 by (simp add: vector-covector)
also have ... = rank * S'+
proof -
  have rankT * ?r ≤ S'+ * rankT * ?s
    using 9 comp-associative by auto
  hence ?r ≤ rank * S'+ * rankT * ?s
    using 4 shunt-mapping comp-associative rank-property-def by auto
  hence ?r * ?sT ≤ rank * S'+ * rankT
    using 5 shunt-bijective by blast
  hence ?r * ?sT * rank ≤ rank * S'+
    using 4 shunt-bijective rank-property-def mapping-conv-bijective by
auto
  thus ?thesis
    using sup-absorb2 by blast
  qed
finally show (?p3 - 1) * rank ≤ rank * S'+
  .
qed
qed
qed
show ¬ ?rr ≤ S'+ * ?rs → ((?rr = ?rs → union-sets-postcondition ?p4
x y p0 ∧ rank-property ?p4 ?rank) ∧
  (?rr ≠ ?rs → union-sets-postcondition ?p4 x y p0
  ∧ rank-property ?p4 rank))
proof
  assume ¬ ?rr ≤ S'+ * ?rs
  hence ?rr ≤ -(S'+ * ?rs)
    using 6 by (meson point-in-vector-or-complement S'-regular
bijective-regular regular-closed-star regular-mult-closed vector-mult-closed)
  also have ... = -S'+ * ?rs
    using 6 comp-bijective-complement by simp
  finally have ?rs ≤ -S'+T * ?rr
    using 6 by (metis bijective-reverse conv-complement)
  also have ... ≤ S'* * ?rr
    using 7 by (simp add: mult-left-isotone)
  also have ... = S'+ * ?rr □ ?rr

```

```

    using star.circ-loop-fixpoint mult-assoc by auto
  finally have 10: ?rs - ?rr ≤ S'+ * ?rr
    using half-shunting by blast
  show ((?rr = ?rs → union-sets-postcondition ?p4 x y p0 ∧
rank-property ?p4 ?rank) ∧
(?rr ≠ ?rs → union-sets-postcondition ?p4 x y p0 ∧ rank-property
?p4 rank))
  proof
    show ?rr = ?rs → union-sets-postcondition ?p4 x y p0 ∧
rank-property ?p4 ?rank
  proof
    assume 11: ?rr = ?rs
    show union-sets-postcondition ?p4 x y p0 ∧ rank-property ?p4 ?rank
  proof
    show union-sets-postcondition ?p4 x y p0
      using 1 2 3 by (simp add: union-sets-1-swap)
    show rank-property ?p4 ?rank
  proof (unfold rank-property-def, intro conjI)
    show univalent ?rank
      using 4 5 6 by (meson S'-univalent read-injective
update-univalent rank-property-def)
    have card-less-eq (roots ?p2) (-(S'+ * rankT * top))
      using 4 rank-property-def by blast
    from this obtain i where 12: injective i ∧ univalent i ∧ regular i
    ∧ roots ?p2 ≤ i * -(S'+ * rankT * top)
      using card-less-eq-def by blast
    let ?i = (i[?s→i[[i * ?rr]])[i * ?rr→i[[?s]])
    have 13: ?i = (i * ?rr □ ?sT * i) ⊔ (-(i * ?rr) □ ?s □ ?rrT * iT
* i) ⊔ (-(i * ?rr) □ -?s □ i)
      by (smt (z3) conv-dist-comp conv-involutive
inf.sup-monoid.add-assoc inf-sup-distrib1 sup-assoc)
    have 14: injective ?i
      apply (rule update-injective-swap)
      subgoal using 12 by simp
      subgoal using 5 by simp
      subgoal using 6 12 injective-mult-closed by simp
      subgoal using 5 comp-associative by simp
    done
    have 15: univalent ?i
      apply (rule update-univalent-swap)
      subgoal using 12 by simp
      subgoal using 5 by simp
      subgoal using 5 by simp
      subgoal using 6 12 injective-mult-closed by simp
      subgoal using 5 comp-associative by simp
    done
    have 16: regular ?i
      using 5 6 12 by (smt (z3) bijective-regular p-dist-inf p-dist-sup
pp-dist-comp regular-closed-inf regular-conv-closed)

```

```

have 17: regular (i * ?rr)
  using 6 12 bijective-regular regular-mult-closed by blast
have 18: find-set-precondition ?p1 y
  using 2 3 find-set-precondition-def
path-compression-postcondition-def path-compression-precondition-def by blast
hence ?s = root ?p1 y
  by (meson find-set-function find-set-postcondition-def)
also have ... = root ?p2 y
  using 3 18 by (smt (z3) find-set-precondition-def
path-compression-postcondition-def path-compression-precondition-def same-root)
also have ... ≤ roots ?p2
  by simp
also have ... ≤ i * -(S+ * rankT * top)
  using 12 by simp
finally have 19: ?s ≤ i * -(S+ * rankT * top)
.
have roots ?p4 ≤ ?i * -(S+ * ?rankT * top)
proof -
  have ?r ≤ -?s
    using 5 8 by (meson order.antisym bijective-regular
point-in-vector-or-complement point-in-vector-or-complement-2)
  hence ?s ⊓ ?rT ⊓ 1 = bot
    by (metis (full-types) bot-least inf.left-commute
inf.sup-monoid.add-commute one-inf-conv pseudo-complement)
  hence ?p4 ⊓ 1 ≤ -?s ⊓ ?p2
    by (smt (z3) bot-least comp-inf.semiring.distrib-left
inf.cobounded2 inf.sup-monoid.add-commute le-supI)
  hence roots ?p4 ≤ roots (-?s ⊓ ?p2)
    by (simp add: mult-left-isotone)
  also have ... = -?s ⊓ roots ?p2
    using 5 inf-assoc vector-complement-closed vector-inf-comp by
auto
  also have ... = (i * ?rr ⊓ -?s ⊓ roots ?p2) ⊔ (-(i * ?rr) ⊓ -?s
⊓ roots ?p2)
    using 17 by (smt (z3) comp-inf.star-plus inf-sup-distrib2
inf-top.right-neutral regular-complement-top)
  also have ... ≤ ?i * -(S+ * ?rankT * top)
    proof (rule sup-least)
      have ?rankT * top = (?r ⊓ (ST * ?rs)T)T * top ⊔ (-?r ⊓
rank)T * top
        using conv-dist-sup mult-right-dist-sup by auto
      also have ... = (?rT ⊓ ST * ?rs) * top ⊔ (-?rT ⊓ rankT) * top
        using conv-complement conv-dist-inf conv-involutive by auto
      also have ... = ST * ?rs * (?r ⊓ top) ⊔ (-?rT ⊓ rankT) * top
        using 5 by (smt (z3) covector-inf-comp-3 inf-commute)
      also have ... = ST * ?rs * (?r ⊓ top) ⊔ rankT * (-?r ⊓ top)
        using 5 by (smt (z3) conv-complement
vector-complement-closed covector-inf-comp-3 inf-commute)
      also have ... = ST * ?rs * ?r ⊔ rankT * -?r

```

by *simp*
 also have $\dots \leq S^{T^*} * ?rs * ?r \sqcup \text{rank}^T * \text{top}$
 using *mult-right-isotone sup-right-isotone* by *force*
 also have $\dots \leq S^{T^*} * ?rs \sqcup \text{rank}^T * \text{top}$
 using *5 6* by (*metis inf.eq-refl mult-assoc*)
 finally have $S^{'+} * ?\text{rank}^T * \text{top} \leq S^{'+} * S^{T^*} * ?rs \sqcup S^{'+} * \text{rank}^T * \text{top}$
 by (*smt comp-associative mult-left-dist-sup mult-right-isotone*)
 also have $\dots = S^{'+} * (S' * S^{T^*}) * ?rs \sqcup S^{'+} * \text{rank}^T * \text{top}$
 by (*smt star-plus mult-assoc*)
 also have $\dots \leq S^{'+} * ?rs \sqcup S^{'+} * \text{rank}^T * \text{top}$
 by (*metis S'-injective comp-right-one mult-left-isotone*)
 also have $\dots = ?rs \sqcup S^{'+} * ?rs \sqcup S^{'+} * \text{rank}^T * \text{top}$
 using *comp-associative star.circ-loop-fixpoint sup-commute* by
 finally have $\dots = ?rs \sqcup S^{'+} * \text{rank}^T * \text{top}$
 by (*smt (verit, del-insts) comp-associative mult-left-dist-sup*
sup.orderE sup-assoc sup-commute top.extremum)
 finally have *20*: $S^{'+} * ?\text{rank}^T * \text{top} \leq ?rs \sqcup S^{'+} * \text{rank}^T * \text{top}$
 .
 have $?s * ?s^T = (?s \sqcap i * -(S^{'+} * \text{rank}^T * \text{top})) * ?s^T$
 using *19 inf.orderE* by *fastforce*
 also have $\dots = (?s \sqcap i * -(S^{'+} * \text{rank}^T * \text{top})) * \text{top} \sqcap ?s^T$
 using *5* by (*smt (z3) covector-comp-inf vector-conv-covector*
vector-covector vector-top-closed)
 also have $\dots = ?s \sqcap i * -(S^{'+} * \text{rank}^T * \text{top}) * \text{top} \sqcap ?s^T$
 using *5 vector-inf-comp* by *auto*
 also have $\dots \leq 1 \sqcap i * -(S^{'+} * \text{rank}^T * \text{top}) * \text{top}$
 using *5* by (*smt (verit, ccfv-SIG) inf.cobounded1*
inf.cobounded2 inf-greatest order-trans vector-covector)
 also have $\dots = 1 \sqcap i * -(S^{'+} * \text{rank}^T * \text{top})$
 using *comp-associative vector-complement-closed*
 also have $\dots \leq 1 \sqcap i * -(S^{'+} * \text{rank}^T)$
 by (*meson comp-inf.mult-right-isotone mult-right-isotone*
p-antitone top-right-mult-increasing)
 also have $\dots \leq 1 \sqcap i * S^{'+} * \text{rank}^T$
 proof –
 have $S^{'+} * \text{rank}^T \sqcup S^{'+} * \text{rank}^T = (S^{T^*} \sqcup S^{'+}) * \text{rank}^T$
 by (*simp add: conv-star-commute mult-right-dist-sup*)
 also have $\dots = (S^{T^*} \sqcup S^{'+}) * \text{rank}^T$
 by (*smt (z3) comp-associative semiring.distrib-right*
star.circ-loop-fixpoint sup.left-commute sup-commute sup-idem)
 also have $\dots = \text{top} * \text{rank}^T$
 by (*simp add: S'-star-connex sup-commute*)
 also have $\dots = \text{top}$
 using *4 rank-property-def total-conv-surjective* by *blast*
 finally have $-(S^{'+} * \text{rank}^T) \leq S^{'+} * \text{rank}^T$

by (*metis half-shunting inf.idem top-greatest*)
 thus *?thesis*
 using *comp-associative inf.sup-right-isotone*
mult-right-isotone by *auto*
 qed
 also have $\dots = 1 \sqcap \text{rank} * S'^* * i^T$
 by (*metis comp-associative conv-dist-comp conv-involutive*
one-inf-conv)
 also have $\dots \leq \text{rank} * S'^* * i^T$
 by *simp*
 finally have $?s \leq \text{rank} * S'^* * i^T * ?s$
 using *5 shunt-bijective* by *auto*
 hence $?rs \leq S'^* * i^T * ?s$
 using *4 shunt-mapping comp-associative rank-property-def* by
auto
 hence $?s * (i * ?rr \sqcap -?s \sqcap \text{roots } ?p2) \leq ?s * (i * S'^* * i^T * ?s \sqcap -?s \sqcap \text{roots } ?p2)$
 using *11 comp-associative comp-inf.mult-left-isotone*
comp-isotone inf.eq-refl by *auto*
 also have $\dots = ?s * ((i * S'^+ * i^T * ?s \sqcup i * i^T * ?s) \sqcap -?s \sqcap \text{roots } ?p2)$
 by (*metis comp-associative mult-left-dist-sup*
star.circ-loop-fixpoint)
 also have $\dots \leq ?s * ((i * S'^+ * i^T * ?s \sqcup 1 * ?s) \sqcap -?s \sqcap \text{roots } ?p2)$
 using *12* by (*metis mult-left-isotone sup-right-isotone*
comp-inf.mult-left-isotone mult-right-isotone)
 also have $\dots = ?s * (i * S'^+ * i^T * ?s \sqcap -?s \sqcap \text{roots } ?p2)$
 using *comp-inf.comp-right-dist-sup* by *simp*
 also have $\dots \leq ?s * (i * S'^+ * i^T * ?s \sqcap \text{roots } ?p2)$
 using *comp-inf.mult-left-isotone inf.cobounded1*
mult-right-isotone by *blast*
 also have $\dots \leq ?s * (i * S'^+ * i^T * ?s \sqcap i * -(S'^+ * \text{rank}^T * \text{top}))$
 using *12 comp-inf.mult-right-isotone mult-right-isotone* by *auto*
 also have $\dots = ?s * (i * S'^+ * i^T * ?s \sqcap i) * -(S'^+ * \text{rank}^T * \text{top})$
 using *5* by (*simp add: comp-associative vector-inf-comp*)
 also have $\dots = (?s \sqcap (i * S'^+ * i^T * ?s)^T) * i * -(S'^+ * \text{rank}^T * \text{top})$
 using *5 covector-inf-comp-3 mult-assoc* by *auto*
 also have $\dots = (?s \sqcap ?s^T * i * S'^{+T} * i^T) * i * -(S'^+ * \text{rank}^T * \text{top})$
 using *conv-dist-comp conv-involutive mult-assoc* by *auto*
 also have $\dots = (?s \sqcap ?s^T) * i * S'^{+T} * i^T * i * -(S'^+ * \text{rank}^T * \text{top})$
 using *5 vector-inf-comp* by *auto*
 also have $\dots \leq i * S'^{+T} * i^T * i * -(S'^+ * \text{rank}^T * \text{top})$
 using *5* by (*metis point-antisymmetric mult-left-isotone*)

mult-left-one)
also have $\dots \leq i * S'^{+T} * -(S'^{+} * rank^T * top)$
using 12 **by** (*smt mult-left-isotone mult-right-isotone*)
mult-assoc comp-right-one)
also have $\dots \leq i * -(S'^{*} * rank^T * top)$
proof –
have $S'^{+} * S'^{*} * rank^T * top \leq S'^{+} * rank^T * top$
by (*simp add: comp-associative star.circ-transitive-equal*)
hence $S'^{+T} * -(S'^{+} * rank^T * top) \leq -(S'^{*} * rank^T * top)$
by (*smt (verit, ccfv-SIG) comp-associative*
conv-complement-sub-leq mult-right-isotone order.trans p-antitone)
thus *?thesis*
by (*simp add: comp-associative mult-right-isotone*)
qed
also have $\dots \leq i * -(S'^{+} * ?rank^T * top)$
proof –
have $S'^{+} * ?rank^T * top \leq ?rs \sqcup S'^{+} * rank^T * top$
using 20 **by** *simp*
also have $\dots \leq rank^T * top \sqcup S'^{+} * rank^T * top$
using *mult-right-isotone sup-left-isotone top.extremum* **by**
blast
also have $\dots = S'^{*} * rank^T * top$
using *comp-associative star.circ-loop-fixpoint sup-commute*
by *auto*
finally show *?thesis*
using *mult-right-isotone p-antitone* **by** *blast*
qed
finally have $?s * (i * ?rr \sqcap -?s \sqcap roots ?p2) \leq i * -(S'^{+} *$
 $?rank^T * top)$
 \cdot
hence $i * ?rr \sqcap -?s \sqcap roots ?p2 \leq ?s^T * i * -(S'^{+} * ?rank^T *$
 $top)$
using 5 *shunt-mapping bijective-conv-mapping mult-assoc* **by**
auto
hence $i * ?rr \sqcap -?s \sqcap roots ?p2 \leq i * ?rr \sqcap ?s^T * i * -(S'^{+} *$
 $?rank^T * top)$
by (*simp add: inf.sup-monoid.add-assoc*)
also have $\dots = (i * ?rr \sqcap ?s^T * i) * -(S'^{+} * ?rank^T * top)$
using 6 *vector-inf-comp vector-mult-closed* **by** *simp*
also have $\dots \leq ?i * -(S'^{+} * ?rank^T * top)$
using 13 *comp-left-increasing-sup sup-assoc* **by** *auto*
finally show $i * ?rr \sqcap -?s \sqcap roots ?p2 \leq ?i * -(S'^{+} * ?rank^T$
 $* top)$
 \cdot
have $-(i * ?rr) \sqcap roots ?p2 \leq -(i * ?rr) \sqcap i * -(S'^{+} * rank^T$
 $* top)$
using 12 *inf.sup-right-isotone* **by** *auto*
also have $\dots \leq -(i * ?rr) \sqcap i * -(?rs \sqcup S'^{+} * rank^T * top)$
proof –

have 21 : *regular* ($?rs \sqcup S'^+ * rank^T * top$)
using $4\ 6$ *rank-property-def mapping-regular S'-regular*
pp-dist-star regular-conv-closed regular-mult-closed bijective-regular
regular-closed-sup **by** *auto*
have $?rs \sqcup S'^+ * rank^T * top \leq S'^+ * rank^T * top \sqcup ?rr$
using 11 **by** *simp*
hence $(?rs \sqcup S'^+ * rank^T * top) - S'^+ * rank^T * top \leq ?rr$
using *half-shunting sup-commute* **by** *auto*
hence $-(S'^+ * rank^T * top) \leq -(?rs \sqcup S'^+ * rank^T * top) \sqcup$
 $?rr$
using 21 **by** (*metis inf.sup-monoid.add-commute*
shunting-var-p sup-commute)
hence $i * -(S'^+ * rank^T * top) \leq i * -(?rs \sqcup S'^+ * rank^T * top) \sqcup i * ?rr$
by (*metis mult-left-dist-sup mult-right-isotone*)
hence $-(i * ?rr) \sqcap i * -(S'^+ * rank^T * top) \leq i * -(?rs \sqcup S'^+ * rank^T * top)$
using *half-shunting inf.sup-monoid.add-commute* **by** *fastforce*
thus *?thesis*
using *inf.le-sup-iff* **by** *blast*
qed
also have $... \leq -(i * ?rr) \sqcap i * -(S'^+ * ?rank^T * top)$
using 20 **by** (*meson comp-inf.mult-right-isotone*
mult-right-isotone p-antitone)
finally have $-(i * ?rr) \sqcap -?s \sqcap roots\ ?p2 \leq -(i * ?rr) \sqcap -?s$
 $\sqcap i * -(S'^+ * ?rank^T * top)$
by (*smt (z3) inf.boundedI inf.cobounded1 inf.coboundedI2*
inf.sup-monoid.add-assoc inf.sup-monoid.add-commute)
also have $... \leq ?i * -(S'^+ * ?rank^T * top)$
using $5\ 6\ 13$ **by** (*smt (z3) sup-commute*
vector-complement-closed vector-inf-comp vector-mult-closed
comp-left-increasing-sup)
finally show $-(i * ?rr) \sqcap -?s \sqcap roots\ ?p2 \leq ?i * -(S'^+ * ?rank^T * top)$
 $?$
qed
finally show *?thesis*
 $?$
qed
thus *card-less-eq* ($roots\ ?p4$) $(-(S'^+ * ?rank^T * top))$
using $14\ 15\ 16$ *card-less-eq-def* **by** *auto*
have $?s \leq i * -(S'^+ * rank^T * top)$
using 19 **by** *simp*
also have $... \leq i * -(S'^+ * ?rr)$
using *mult-right-isotone p-antitone top.extremum mult-assoc* **by**
auto
also have $... = i * -S'^+ * ?rr$
using 6 *comp-bijective-complement mult-assoc* **by** *fastforce*
finally have $?rr \leq -S'^{T+} * i^T * ?s$

using 5 6 **by** (*metis conv-complement conv-dist-comp conv-plus-commute bijective-reverse*)
also have $\dots \leq S'^* * i^T * ?s$
using 7 *conv-plus-commute mult-left-isotone* **by** *auto*
finally have 22: $?rr \leq S'^* * i^T * ?s$
.

have $?r = \text{root } ?p1 \ x$
using 2 *path-compression-postcondition-def*
path-compression-precondition-def **by** *blast*
also have $\dots = \text{root } ?p2 \ x$
using 3 18 **by** (*smt (z3) find-set-precondition-def path-compression-postcondition-def path-compression-precondition-def same-root*)
also have $\dots \leq \text{roots } ?p2$
by *simp*
also have $\dots \leq i * -(S'^+ * \text{rank}^T * \text{top})$
using 12 **by** *simp*
also have $\dots \leq i * -(S'^+ * ?rr)$
using *mult-right-isotone p-antitone top.extremum mult-assoc* **by** *auto*

also have $\dots = i * -S'^+ * ?rr$
using 6 *comp-bijective-complement mult-assoc* **by** *fastforce*
finally have $?rr \leq -S'^+ * i^T * ?r$
using 5 6 **by** (*metis conv-complement conv-dist-comp conv-plus-commute bijective-reverse*)
also have $\dots \leq S'^* * i^T * ?r$
using 7 *conv-plus-commute mult-left-isotone* **by** *auto*
finally have $?rr \leq S'^* * i^T * ?r$
.

hence $?rr \leq S'^* * i^T * ?r \sqcap S'^* * i^T * ?s$
using 22 *inf.boundedI* **by** *blast*
also have $\dots = (S'^+ * i^T * ?r \sqcup i^T * ?r) \sqcap S'^* * i^T * ?s$
by (*simp add: star.circ-loop-fixpoint mult-assoc*)
also have $\dots \leq S'^+ * i^T * ?r \sqcup (i^T * ?r \sqcap S'^* * i^T * ?s)$
by (*metis comp-inf.mult-right-dist-sup eq-refl inf.cobounded1 semiring.add-mono*)
also have $\dots \leq S' * \text{top} \sqcup (i^T * ?r \sqcap S'^* * i^T * ?s)$
using *comp-associative mult-right-isotone sup-left-isotone top.extremum* **by** *auto*
also have $\dots = S' * \text{top} \sqcup (i^T * ?r \sqcap (S'^+ * i^T * ?s \sqcup i^T * ?s))$
by (*simp add: star.circ-loop-fixpoint mult-assoc*)
also have $\dots \leq S' * \text{top} \sqcup S'^+ * i^T * ?s \sqcup (i^T * ?r \sqcap i^T * ?s)$
by (*smt (z3) comp-inf.semiring.distrib-left inf.sup-right-divisibility star.circ-loop-fixpoint sup-assoc sup-commute sup-inf-distrib1*)
also have $\dots \leq S' * \text{top} \sqcup (i^T * ?r \sqcap i^T * ?s)$
by (*metis comp-associative mult-right-isotone order.refl sup.orderE top.extremum*)
also have $\dots = S' * \text{top} \sqcup i^T * (?r \sqcap ?s)$
using 12 *conv-involutive univalent-comp-left-dist-inf* **by** *auto*
also have $\dots = S' * \text{top}$

using 5 8 *distinct-points* **by** *auto*
finally have $top \leq ?rr^T * S' * top$
using 6 **by** (*smt conv-involutive shunt-mapping*
bijjective-conv-mapping mult-assoc)
hence surjective ($S'^T * ?rs$)
using 6 11 **by** (*smt conv-dist-comp conv-involutive*
point-conv-comp top-le)
thus total $?rank$
using 4 5 *bijjective-regular update-total rank-property-def* **by** *blast*
show $(?p4 - 1) * ?rank \leq ?rank * S'^+$
proof –
have 23: *univalent ?p2*
using 3 *path-compression-postcondition-def*
path-compression-precondition-def **by** *blast*
have 24: $?r \sqcap (?p4 - 1) * ?rank \leq ?s^T * rank * S' * S'^+$
proof –
have $?r \sqcap (?p4 - 1) * ?rank = (?r \sqcap ?p4 \sqcap -1) * ?rank$
using 5 *vector-complement-closed vector-inf-comp inf-assoc* **by**
auto
also have $... = (?r \sqcap -?s \sqcap ?p4 \sqcap -1) * ?rank$
using 5 8 **by** (*smt (z3) order.antisym bijjective-regular*
point-in-vector-or-complement point-in-vector-or-complement-2 inf-absorb1)
also have $... = (?r \sqcap -?s \sqcap ?p2 \sqcap -1) * ?rank$
by (*simp add: inf.left-commute inf.sup-monoid.add-commute*
inf-sup-distrib1 inf-assoc)
also have $... \leq (?r \sqcap ?p2 \sqcap -1) * ?rank$
using *inf.sup-left-isotone inf-le1 mult-left-isotone* **by** *blast*
also have $... = bot$
proof –
have $?r = root\ ?p1\ x$
using 2 *path-compression-postcondition-def*
path-compression-precondition-def **by** *blast*
also have $... = root\ ?p2\ x$
using 3 18 **by** (*smt (z3) find-set-precondition-def*
path-compression-postcondition-def path-compression-precondition-def same-root)
also have $... \leq roots\ ?p2$
by *simp*
finally have $?r \sqcap ?p2 \leq roots\ ?p2 \sqcap ?p2$
using *inf.sup-left-isotone* **by** *blast*
also have $... \leq (?p2 \sqcap 1) * (?p2 \sqcap 1)^T * ?p2$
by (*smt (z3) comp-associative comp-inf.star-plus dedekind-1*
inf-top-right order-lesseq-imp)
also have $... = (?p2 \sqcap 1) * (?p2 \sqcap 1) * ?p2$
using *coreflexive-symmetric* **by** *force*
also have $... \leq (?p2 \sqcap 1) * ?p2$
by (*metis coreflexive-comp-top-inf inf.cobounded2*
mult-left-isotone)
also have $... \leq ?p2 \sqcap 1$
by (*smt 23 comp-inf.mult-semi-associative conv-dist-comp*

conv-dist-inf conv-order equivalence-one-closed inf.absorb1
inf.sup-monoid.add-assoc injective-codomain)
also have $\dots \leq 1$
by *simp*
finally have $?r \sqcap ?p2 \leq 1$
 \cdot
thus *?thesis*
by (*metis pseudo-complement regular-one-closed*
semiring.mult-not-zero)
qed
finally show *?thesis*
using *bot-least le-bot by blast*
qed
have $?5: -?r \sqcap (?p4 - 1) * ?rank \leq rank * S'^+$
proof $-$
have $-?r \sqcap (?p4 - 1) * ?rank = (-?r \sqcap ?p4 \sqcap -1) * ?rank$
using *5 vector-complement-closed vector-inf-comp inf-assoc by*
auto
also have $\dots = (-?r \sqcap (?s \sqcup -?s) \sqcap ?p4 \sqcap -1) * ?rank$
using *5 bijective-regular inf-top-right regular-complement-top*
by *auto*
also have $\dots = (-?r \sqcap ?s \sqcap ?p4 \sqcap -1) * ?rank \sqcup (-?r \sqcap -?s$
 $\sqcap ?p4 \sqcap -1) * ?rank$
by (*smt (z3) inf-sup-distrib1 inf-sup-distrib2*
mult-right-dist-sup)
also have $\dots = (-?r \sqcap ?s \sqcap ?r^T \sqcap -1) * ?rank \sqcup (-?r \sqcap -?s$
 $\sqcap ?p2 \sqcap -1) * ?rank$
using *5 by (smt (z3) bijective-regular*
comp-inf.comp-left-dist-sup inf-assoc inf-commute inf-top-right mult-right-dist-sup
regular-complement-top)
also have $\dots \leq (?s \sqcap ?r^T \sqcap -1) * ?rank \sqcup (-?s \sqcap ?p2 \sqcap -1)$
 $* ?rank$
by (*smt (z3) comp-inf.semiring.distrib-left inf.cobounded2*
inf.sup-monoid.add-assoc mult-left-isotone mult-right-dist-sup)
also have $\dots \leq (?s \sqcap ?r^T) * ?rank \sqcup (?p2 - 1) * ?rank$
by (*smt (z3) inf.cobounded1 inf.cobounded2*
inf.sup-monoid.add-assoc mult-left-isotone semiring.add-mono)
also have $\dots = ?s * (?r \sqcap ?rank) \sqcup (?p2 - 1) * ?rank$
using *5 by (simp add: covector-inf-comp-3)*
also have $\dots = ?s * (?r \sqcap (S'^T * ?rs)^T) \sqcup (?p2 - 1) * ?rank$
using *inf-commute update-inf-same mult-assoc by force*
also have $\dots = ?s * (?r \sqcap ?s^T * rank * S') \sqcup (?p2 - 1) * ?rank$
using *comp-associative conv-dist-comp conv-involutive by auto*
also have $\dots \leq ?s * ?s^T * rank * S' \sqcup (?p2 - 1) * ?rank$
using *comp-associative inf.cobounded2 mult-right-isotone*
semiring.add-right-mono by auto
also have $\dots \leq 1 * rank * S' \sqcup (?p2 - 1) * ?rank$
using *5 by (meson mult-left-isotone order.refl*
semiring.add-mono)

also have ... = rank * S' \sqcup (?p2 - 1) * (?r \sqcap (S^T * ?rr)^T) \sqcup
 (?p2 - 1) * (-?r \sqcap rank)
using 11 comp-associative mult-1-left mult-left-dist-sup
 sup-assoc **by auto**
also have ... \leq rank * S' \sqcup (?p2 - 1) * (?r \sqcap ?r^T * rank * S')
 \sqcup (?p2 - 1) * rank
using comp-associative conv-dist-comp conv-involutive
 inf.cobounded1 inf.sup-monoid.add-commute mult-right-isotone
 semiring.add-left-mono **by auto**
also have ... = rank * S' \sqcup (?p2 - 1) * (?r \sqcap ?r^T) * rank * S'
 \sqcup (?p2 - 1) * rank
using 5 comp-associative vector-inf-comp **by auto**
also have ... \leq rank * S' \sqcup (?p2 - 1) * rank * S' \sqcup (?p2 - 1)
 * rank
using 5 **by** (metis point-antisymmetric mult-left-isotone
 mult-right-isotone sup-left-isotone sup-right-isotone comp-right-one)
also have ... \leq rank * S' \sqcup rank * S⁺ * S' \sqcup (?p2 - 1) * rank
using 4 **by** (metis rank-property-def mult-left-isotone
 sup-left-isotone sup-right-isotone)
also have ... \leq rank * S' \sqcup rank * S⁺ * S' \sqcup rank * S⁺
using 4 **by** (metis rank-property-def sup-right-isotone)
also have ... \leq rank * S⁺
using comp-associative eq-refl le-sup-iff mult-right-isotone
 star.circ-mult-increasing star.circ-plus-same star.left-plus-below-circ **by auto**
finally show ?thesis
 .
qed
have (?p4 - 1) * ?rank = (?r \sqcap (?p4 - 1) * ?rank) \sqcup (-?r \sqcap
 (?p4 - 1) * ?rank)
using 5 **by** (smt (verit, ccfv-threshold) bijective-regular
 inf-commute inf-sup-distrib2 inf-top-right regular-complement-top)
also have ... \leq (?r \sqcap ?s^T * rank * S' * S⁺) \sqcup (-?r \sqcap rank *
 S⁺)
using 24 25 **by** (meson inf.boundedI inf.cobounded1
 semiring.add-mono)
also have ... = (?r \sqcap ?s^T * rank * S') * S⁺ \sqcup (-?r \sqcap rank) *
 S⁺
using 5 vector-complement-closed vector-inf-comp **by auto**
also have ... = ?rank * S⁺
using conv-dist-comp mult-right-dist-sup **by auto**
finally show ?thesis
 .
qed
qed
qed
qed
show ?rr \neq ?rs \longrightarrow union-sets-postcondition ?p4 x y p0 \wedge
 rank-property ?p4 rank
proof

```

assume ?rr ≠ ?rs
hence ?rs ⊓ ?rr = bot
using 6 by (meson bijective-regular dual-order.eq-iff
point-in-vector-or-complement point-in-vector-or-complement-2
pseudo-complement)
hence 26: ?rs ≤ S+ * ?rr
using 10 le-iff-inf pseudo-complement by auto
show union-sets-postcondition ?p4 x y p0 ∧ rank-property ?p4 rank
proof
show union-sets-postcondition ?p4 x y p0
using 1 2 3 by (simp add: union-sets-1-swap)
show rank-property ?p4 rank
proof (unfold rank-property-def, intro conjI)
show univalent rank total rank
using 1 rank-property-def by auto
have ?r ≤ -?s
using 5 8 by (meson order.antisym bijective-regular
point-in-vector-or-complement point-in-vector-or-complement-2)
hence ?s ⊓ ?rT ⊓ 1 = bot
by (metis (full-types) bot-least inf.left-commute
inf.sup-monoid.add-commute one-inf-conv pseudo-complement)
hence ?p4 ⊓ 1 ≤ ?p2
by (smt half-shunting inf.cobounded2 pseudo-complement
regular-one-closed semiring.add-mono sup-commute)
hence roots ?p4 ≤ roots ?p2
by (simp add: mult-left-isotone)
thus card-less-eq (roots ?p4) (-(S+ * rankT * top))
using 4 by (meson rank-property-def card-less-eq-def order-trans)
have (?p4 - 1) * rank = (?s ⊓ ?rT ⊓ -1) * rank ⊔ (-?s ⊓ ?p2
⊓ -1) * rank
using comp-inf.semiring.distrib-right mult-right-dist-sup by auto
also have ... ≤ (?s ⊓ ?rT ⊓ -1) * rank ⊔ (?p2 - 1) * rank
using comp-inf.mult-semi-associative mult-left-isotone
sup-right-isotone by auto
also have ... ≤ (?s ⊓ ?rT ⊓ -1) * rank ⊔ rank * S+
using 4 sup-right-isotone rank-property-def by blast
also have ... ≤ (?s ⊓ ?rT) * rank ⊔ rank * S+
using inf-le1 mult-left-isotone sup-left-isotone by blast
also have ... = ?s * ?rT * rank ⊔ rank * S+
using 5 by (simp add: vector-covector)
also have ... = rank * S+
proof -
have rankT * ?s ≤ S+ * rankT * ?r
using 26 comp-associative by auto
hence ?s ≤ rank * S+ * rankT * ?r
using 4 shunt-mapping comp-associative rank-property-def by
auto
hence ?s * ?rT ≤ rank * S+ * rankT
using 5 shunt-bijective by blast

```

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    hence ?s * ?rT * rank ≤ rank * S+
    using 4 shunt-bijective rank-property-def mapping-conv-bijective
by auto
    thus ?thesis
    using sup-absorb2 by blast
qed
finally show (?p4 - 1) * rank ≤ rank * S+
.
qed
qed
qed
qed
qed
qed
qed
show ?r = ?s → union-sets-postcondition ?p5 x y p0 ∧ rank-property ?p5
rank
proof
  assume 27: ?r = ?s
  show union-sets-postcondition ?p5 x y p0 ∧ rank-property ?p5 rank
  proof
    show union-sets-postcondition ?p5 x y p0
    using 1 2 3 27 by (simp add: union-sets-1-skip)
    show rank-property ?p5 rank
    using 4 27 by simp
  qed
qed
qed
qed
end
end

```

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