Ramsey's Theorem

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Abstract

The infinite form of Ramsey's Theorem is proved following Boolos and Jeffrey, Chapter 26.

Contents

1 Ramsey's Theorem

theory *Ramsey*
imports *Main HOL−Library.Infinite-Set*
begin

declare [[*simp-depth-limit* = *5*]]

1.1 Library lemmas

lemma *infinite-inj-infinite-image*: *infinite Z* ⟹ *inj-on f Z* ⟹ *infinite* (*f* ‘ *Z*)
apply(*rule ccontr*)
apply(*simp*)
apply(*force dest*: *finite-imageD*)
done

lemma *infinite-dom-finite-rng*: [| *infinite A*; *finite* (*f* ‘ *A*) |] ==> ∃ *b* ∈ *f* ‘ *A*. *infinite* {*a* : *A*. *f a* = *b*}
apply(*rule ccontr*) **apply**(*simp*)
apply(*subgoal-tac* ⋃((λ*b*. {*a* : *A*. *f a* = *f b*}) ‘ *A*) = *A*) **prefer** *2* **apply**(*blast*)
apply(*subgoal-tac* (*UN c* : *f* ‘ *A*. {*a* : *A*. *f a* = *c*})= (*UN b*:*A*. {*a* : *A*. *f a* = *f b*})) **prefer** *2* **apply**(*blast*)
apply(*subgoal-tac finite* (*UN c*:*f* ‘ *A*. {*a* : *A*. *f a* = *c*})) **apply**(*force*)

apply(*rule finite-UN-I*)
apply(*auto*)
done

lemma *infinite-mem*: *infinite* $X \Longrightarrow \exists\, x.\ x \in X$
apply(*rule ccontr*)
apply(*force*)
done

lemma *not-empty-least*: $(Y{::}nat\ set) \neq \{\} \Longrightarrow \exists\, m.\ m \in Y \wedge (\forall\, m'.\ m' \in Y \longrightarrow m \leq m')$
apply(*rule-tac x=LEAST x. x : Y* **in** *exI*)
apply(*rule*)
apply(*rule LeastI-ex*) **apply**(*force*)
apply(*rule*)
apply(*rule*)
apply(*rule Least-le*)
apply(*assumption*)
done

1.2 Dependent Choice Variant

—
primrec *choice* :: $('a \Rightarrow bool) \Rightarrow ('a \Rightarrow 'a \Rightarrow bool) \Rightarrow nat \Rightarrow 'a$ **where**
choice P R 0 = (*SOME x. P x*)
| *choice P R* (*Suc n*) = (*let x = choice P R n in SOME y. P y* $\wedge$ *R x y*)
—

lemma *dc*:
$(\forall\, x\ y\ z.\ R\ x\ y \wedge R\ y\ z \longrightarrow R\ x\ z)$
$\wedge\ (\exists\, x0.\ P\ x0)$
$\wedge\ (\forall\, x.\ P\ x \longrightarrow (\exists\, y.\ P\ y \wedge R\ x\ y))$
$\longrightarrow\ (\exists f{::}nat \Rightarrow 'b.\ (\forall\, n.\ P\ (f\ n)) \wedge (\forall\, n\ m.\ R\ (f\ n)\ (f\ (n+m+1))))$

apply(*intro allI impI*, *elim exE conjE*)
apply(*rule-tac x=choice P R* **in** *exI*)
apply(*subgoal-tac* $(\forall\, n.\ P\ (choice\ P\ R\ n))$) **prefer** *2*
apply(*rule*, *induct-tac n*)
apply(*simp add*: *Let-def*) **apply**(*rule someI-ex*) **apply**(*blast*)
apply(*simp add*: *Let-def*) **apply**(*subgoal-tac P* (*SOME y. P y* $\wedge$ *R* (*choice P R na*) *y*) $\wedge$ *R* (*choice P R na*) (*SOME y. P y* $\wedge$ *R* (*choice P R na*) *y*)) **apply**(*blast*)
apply(*rule someI-ex*) **apply**(*blast*)
apply(*rule*) **apply**(*assumption*)

apply(*subgoal-tac* $\forall\, n.\ R\ (choice\ P\ R\ n)\ (choice\ P\ R\ (Suc\ n))$) **prefer** *2*
apply(*rule*)
apply(*simp add*: *Let-def*)
apply(*subgoal-tac P* (*SOME y. P y* $\wedge$ *R* (*choice P R n*) *y*) $\wedge$ *R* (*choice P R n*) (*SOME y. P y* $\wedge$ *R* (*choice P R n*) *y*)) **apply**(*blast*)

apply(*rule someI-ex*) **apply**(*force*)

apply(*rule*) **apply**(*rule*)
apply(*induct-tac m*)
apply(*force*)
apply(*drule-tac x=n+na+1* **in** *spec*) **back**
apply(*force simp del*: *choice.simps*)
done

1.3 Partitions

definition
part :: *nat* ⇒ *nat* ⇒ *'a set* ⇒ (*'a set* ⇒ *nat*) ⇒ *bool* **where**
part r s Y f = (∀ *X*. *X* ⊆ *Y* ∧ *finite X* ∧ *card X* = *r* ⟶ *f X* < *s*)

lemma *part*: [| *infinite YY*; *part* (*Suc n*) *s YY f*; *yy* : *YY* |] ==> *part n s* (*YY* − {*yy*}) (λ*u*. *f* (*insert yy u*))
apply(*simp add*: *part-def*)
apply(*intro allI impI*, *elim exE conjE*)
apply(*drule-tac x=insert yy X* **in** *spec*)
apply(*force simp*:*card-Diff-singleton-if*)
done

lemma *part-subset*: *part* (*Suc n*) *s YY f* ⟹ *Y* ⊆ *YY* ⟹ *part* (*Suc n*) *s Y f*
apply(*simp add*: *part-def*)
apply(*blast*)
done

1.4 Ramsey's theorem

lemma *ramsey*:
∀(*s*::*nat*) (*r*::*nat*) (*YY*::*'a set*) (*f*::*'a set* ⇒ *nat*).
infinite YY
∧ (∀ *X*. *X* ⊆ *YY* ∧ *finite X* ∧ *card X* = *r* ⟶ *f X* < *s*)
⟶ (∃ *Y' t'*.
Y' ⊆ *YY*
∧ *infinite Y'*
∧ *t'* < *s*
∧ (∀ *X*. *X* ⊆ *Y'* ∧ *finite X* ∧ *card X* = *r* ⟶ *f X* = *t'*))
apply(*simp add*: *part-def*[*symmetric*] *del*: *ex-simps*)
apply(*rule*, *rule*, *rule-tac nat.induct*)

apply(*intro allI impI*)
apply(*rule-tac x=YY* **in** *exI*)
apply(*rule-tac x=f* {} **in** *exI*)
apply(*force simp*: *part-def*)

apply(*intro allI impI*) **apply**(*elim exE conjE*)

apply(*subgoal-tac*
$\exists\, g.$
$(\forall\, m{::}nat.\ let\ (y,Y,t) = (g\ m)\ in$
$y \in YY \wedge y \notin Y$
$\wedge\ Y \subseteq YY \wedge infinite\ Y$
$\wedge\ t < s$
$\wedge\ (\forall\, X.\ X \subseteq Y \wedge finite\ X \wedge card\ X = nat \longrightarrow (f \circ insert\ y)\ X = t))$
$\wedge\ (\forall\, m\ m'.$
$let\ (y,Y,t) = (g\ m)\ in$
$let\ (y',Y',t') = (g\ (m{+}m'{+}1))\ in$
$y' : Y$
$\wedge\ Y' \subseteq Y$
)
)
prefer *2*
apply(*cut-tac*

$P = \lambda gn.$
$let\ (y,Y,t) = (gn)\ in$
$y \in YY \wedge y \notin Y$
$\wedge\ Y \subseteq YY \wedge infinite\ Y$
$\wedge\ t < s$
$\wedge\ (\forall\, X.\ X \subseteq Y \wedge finite\ X \wedge card\ X = nat \longrightarrow (f \circ insert\ y)\ X = t)$

and $R = \lambda gn\ gn'.$
$let\ (y,Y,t) = (gn)\ in$
$let\ (y',Y',t') = (gn')\ in$
$y' : Y$
$\wedge\ Y' \subseteq Y$

in *dc*)
apply(*erule impE*)

apply(*intro conjI*)
apply(*intro allI impI*, *elim conjE*) **apply**(*simp add*: *Let-def split-beta*) **apply**(*blast*)
apply(*subgoal-tac* $\exists\, yy.\ yy \in YY$) **prefer** *2* **apply**(*rule infinite-mem*) **apply**(*assumption*)
apply(*elim exE conjE*)
apply(*drule-tac* $x{=}YY - \{yy\}$ **in** *spec*) **apply**(*drule-tac* $x{=}f \circ insert\ yy$ **in** *spec*)
apply(*erule impE*) **apply**(*simp*) **apply**(*rule part*) **apply**(*assumption*+)
apply(*elim exE conjE*)
apply(*rule-tac* $x{=}(yy,Y',t')$ **in** *exI*) **apply**(*simp*) **apply**(*blast*)
apply(*intro allI impI*)
apply(*case-tac x*) **apply**(*rename-tac yx b Yx tx*)
apply(*subgoal-tac* $\exists\, yx'.\ yx' \in Yx$) **prefer** *2* **apply**(*rule infinite-mem*) **ap-**

ply(*force*)
apply(*elim exE conjE*)
apply(*drule-tac x=Yx − {yx′} in spec*)
apply(*drule-tac x=f ∘ insert yx′ in spec*)
apply(*erule impE*) **apply**(*simp*) **apply**(*elim exE conjE*) **apply**(*rule part*) **apply**(*assumption+*)
apply(*rule part-subset*) **apply**(*assumption*) **apply**(*assumption*) **apply**(*assumption*)
apply(*elim exE conjE*)
apply(*rule-tac x=(yx′,Y′,t′) in exI*) **apply**(*simp*) **apply**(*blast*)
apply(*assumption*)

apply(*elim exE conjE*)

apply(*subgoal-tac ∃ s′. s′ < s ∧ infinite {n. (λn. let (y,Y,t) = g n in t) n = s′}*) **prefer** *2*

apply(*subgoal-tac ∃ s′ ∈ ((λn. let (Y, y, t) = g n in t) ‘ UNIV). infinite {n : UNIV. (let (Y, y, t) = g n in t) = s′}*) **prefer** *2*
apply(*rule infinite-dom-finite-rng*) **apply**(*simp*)
apply(*simp (no-asm) add: finite-nat-iff-bounded*)
apply(*rule-tac x=s in exI*)
apply(*rule*)
apply(*simp add: image-iff*) **apply**(*elim exE conjE*)
apply(*drule-tac x=xa in spec*) **apply**(*force simp add: Let-def split-beta*)
apply(*elim bexE conjE*) **apply**(*rule-tac x=s′ in exI*) **apply**(*simp*)
apply(*simp add: image-iff*) **apply**(*elim exE conjE*) **apply**(*drule-tac x=x in spec*) **apply**(*force simp: Let-def split-beta*)

apply(*elim exE conjE*)
apply(*rule-tac x=(λn. let (y,Y,t) = g n in y) ‘ {n. (λn. let (y,Y,t) = g n in t) n = s′} in exI*)
apply(*rule-tac x=s′ in exI*)

apply(*subgoal-tac inj (λn. let (y, Y, t) = g n in y)*) **prefer** *2*
apply(*simp add: inj-on-def*)

apply(*subgoal-tac ∀ x y. x < y ∧ (let (y, Y, t) = g x in y) = (let (y, Y, t) = g y in y) ⟶ x = y*)
apply(*intro allI impI*)
apply(*subgoal-tac x < y | x = y | y < x*) **prefer** *2* **apply**(*arith*)
apply(*elim disjE*)
apply(*drule-tac x=x in spec*) **back back apply**(*drule-tac x=y in spec*) **back back apply**(*force simp: Let-def*)
apply(*assumption*)
apply(*drule-tac x=y in spec*) **back back apply**(*drule-tac x=x in spec*) **back back apply**(*force simp: Let-def*)
apply(*intro allI impI*)
apply(*drule-tac x=x in spec*) **apply**(*drule-tac x=x in spec*) **apply**(*drule-tac x=y−(Suc x) in spec*) **apply**(*force simp: Let-def*)

apply(*intro allI conjI*)

apply(*drule-tac x=xa* **in** *spec*)
apply(*force simp add*: *Let-def split-beta*)

apply(*rule infinite-inj-infinite-image*) **apply**(*assumption*)
apply(*simp add*: *inj-on-def*)

apply(*simp*)

apply(*intro allI impI*, *elim exE conjE*)
apply(*simp add*: *subset-image-iff*) **apply**(*elim exE conjE*)
apply(*subgoal-tac* $\exists\, a.\ a \in AA \wedge (\forall\, a'.\ a' \in AA \longrightarrow a \leq a')$) **prefer** *2* **apply**(*rule not-empty-least*) **apply**(*force*)
apply(*elim exE conjE*)
apply(*case-tac g a rule*: *prod.exhaust*)
apply(*rename-tac b*) **apply**(*case-tac b rule*: *prod.exhaust*) **apply**(*rename-tac ya b Ya ta*)
apply(*subgoal-tac ya* : *X*) **prefer** *2* **apply**(*force intro*!: *rev-image-eqI simp*: *Let-def*)
apply(*drule-tac s=X* **in** *sym*)
apply(*subgoal-tac f X* = (*f* ∘ *insert ya*) (*X* − {*ya*})) **apply**(*simp*)
apply(*drule-tac x=a* **in** *spec*)
apply(*simp add*: *Let-def*) **apply**(*elim exE conjE*)
apply(*drule-tac x=X−*{*ya*} **in** *spec*) **back** **apply**(*erule impE*)
apply(*simp*)
apply(*rule*)
apply(*rule*)
apply(*drule-tac t=X* **in** *sym*) **apply**(*simp*)
apply(*simp add*: *image-iff*) **apply**(*elim bexE exE conjE*) **apply**(*rename-tac a'*)
apply(*subgoal-tac a'* ≠ *a*) **prefer** *2* **apply**(*force*)
apply(*drule-tac x=a* **in** *spec*)
apply(*drule-tac x=a'−Suc a* **in** *spec*) **back**
apply(*simp add*: *Let-def split-beta*) **apply**(*case-tac g a' rule*: *prod.exhaust*) **apply**(*case-tac ba rule*: *prod.exhaust*) **apply**(*rename-tac ya' ba Ya' ta'*) **apply**(*simp add*: *Let-def split-beta*)
apply(*drule-tac x=a'* **in** *spec*) **apply**(*erule impE*) **apply**(*force*)
apply(*force*)
apply(*force simp add*: *card-Diff-singleton-if*)
apply(*subgoal-tac ta* = *s'*) **apply**(*simp*) **apply**(*force*)
apply(*simp*) **apply**(*rule-tac f=f* **in** *arg-cong*) **apply**(*force*)
done

end