

Two theorems about the geometry of the critical points of a complex polynomial

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Abstract

This entry formalises two well-known results about the geometric relation between the *roots* of a complex polynomial and its *critical points*, i.e. the roots of its derivative.

The first of these is the *Gauß–Lucas Theorem*: The critical points of a complex polynomial lie inside the convex hull of its roots.

The second one is *Jensen’s Theorem*: Every non-real critical point of a real polynomial lies inside a disc between two conjugate roots. These discs are called the *Jensen discs*: the Jensen disc of a pair of conjugate roots $a \pm bi$ is the smallest disc that contains both of them, i.e. the disc with centre a and radius b .

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1 Missing Library Material

```
theory Polynomial_Crit_Geometry_Library
```

```
imports
```

```
  "HOL-Computational_Algebra.Computational_Algebra"
```

```
  "HOL-Library.FuncSet"
```

```
  "Polynomial_Interpolation.Ring_Hom_Poly"
```

```
begin
```

1.1 Multisets

```
lemma size_repeat_mset [simp]: "size (repeat_mset n A) = n * size A"  
  by (induction n) auto
```

```
lemma count_image_mset_inj:  
  "inj f  $\implies$  count (image_mset f A) (f x) = count A x"  
  by (induction A) (auto dest!: injD)
```

```
lemma count_le_size: "count A x  $\leq$  size A"  
  by (induction A) auto
```

```
lemma image_mset_cong_simp:  
  " $M = M' \implies (\bigwedge x. x \in\# M \implies f x = g x) \implies \{f x. x \in\# M\} = \{g x. x \in\# M'\}$ "  
  unfolding simp_implies_def by (auto intro: image_mset_cong)
```

```
lemma sum_mset_nonneg:  
  fixes A :: "'a :: ordered_comm_monoid_add multiset"  
  assumes " $\bigwedge x. x \in\# A \implies x \geq 0$ "  
  shows "sum_mset A  $\geq 0$ "  
  using assms by (induction A) auto
```

```
lemma sum_mset_pos:  
  fixes A :: "'a :: ordered_comm_monoid_add multiset"  
  assumes "A  $\neq$  {}"  
  assumes " $\bigwedge x. x \in\# A \implies x > 0$ "  
  shows "sum_mset A  $> 0$ "
```

```
proof -
```

```
  from assms obtain x where "x  $\in\#$  A"
```

```
    by auto
```

```
  hence "A = {x} + (A - {x})"
```

```
    by auto
```

```
  also have "sum_mset ... = x + sum_mset (A - {x})"
```

```
    by simp
```

```
  also have "...  $> 0$ "
```

```
  proof (rule add_pos_nonneg)
```

```
    show "x  $> 0$ "
```

```
      using <x  $\in\#$  A> assms by auto
```

```
    show "sum_mset (A - {x})  $\geq 0$ "
```

```
      using assms sum_mset_nonneg by (metis in_diffD order_less_imp_le)
```

```

qed
finally show ?thesis .
qed

```

1.2 Polynomials

```

lemma order_pos_iff: "p ≠ 0 ⇒ order x p > 0 ↔ poly p x = 0"
  by (cases "order x p = 0") (auto simp: order_root order_0I)

```

```

lemma order_prod_mset:
  "0 ∉# P ⇒ order x (prod_mset P) = sum_mset (image_mset (λp. order
x p) P)"
  by (induction P) (auto simp: order_mult)

```

```

lemma order_prod:
  "(∧x. x ∈ I ⇒ f x ≠ 0) ⇒ order x (prod f I) = (∑ i∈I. order x
(f i))"
  by (induction I rule: infinite_finite_induct) (auto simp: order_mult)

```

```

lemma order_linear_factor:
  assumes "a ≠ 0 ∨ b ≠ 0"
  shows "order x [:a, b:] = (if b * x + a = 0 then 1 else 0)"
proof (cases "b * x + a = 0")
  case True
  have "order x [:a, b:] ≤ degree [:a, b:]"
    using assms by (intro order_degree) auto
  also have "... ≤ 1"
    by simp
  finally have "order x [:a, b:] ≤ 1" .
  moreover have "order x [:a, b:] > 0"
    using assms True by (subst order_pos_iff) (auto simp: algebra_simps)
  ultimately have "order x [:a, b:] = 1"
    by linarith
  with True show ?thesis
    by simp
qed (auto intro!: order_0I simp: algebra_simps)

```

```

lemma order_linear_factor' [simp]:
  assumes "a ≠ 0 ∨ b ≠ 0" "b * x + a = 0"
  shows "order x [:a, b:] = 1"
  using assms by (subst order_linear_factor) auto

```

```

lemma degree_prod_mset_eq: "0 ∉# P ⇒ degree (prod_mset P) = (∑ p∈#P.
degree p)"
  for P :: "'a::idom poly multiset"
  by (induction P) (auto simp: degree_mult_eq)

```

```

lemma degree_prod_list_eq: "0 ∉ set ps ⇒ degree (prod_list ps) = (∑ p←ps.
degree p)"

```

```

for ps :: "'a::idom poly list"
by (induction ps) (auto simp: degree_mult_eq prod_list_zero_iff)

lemma order_conv_multiplicity:
  assumes "p ≠ 0"
  shows "order x p = multiplicity [:-x, 1:] p"
  using assms order[of p x] multiplicity_eqI by metis

1.3 Polynomials over algebraically closed fields

lemma irreducible_alg_closed_imp_degree_1:
  assumes "irreducible (p :: 'a :: alg_closed_field poly)"
  shows "degree p = 1"
proof -
  have "¬(degree p > 1)"
    using assms alg_closed_imp_reducible by blast
  moreover from assms have "degree p ≠ 0"
    by (auto simp: irreducible_def is_unit_iff_degree)
  ultimately show ?thesis
    by linarith
qed

lemma prime_poly_alg_closedE:
  assumes "prime (q :: 'a :: {alg_closed_field, field_gcd} poly)"
  obtains c where "q = [:-c, 1:]" "poly q c = 0"
proof -
  from assms have "degree q = 1"
    by (intro irreducible_alg_closed_imp_degree_1 prime_elem_imp_irreducible)
  auto
  then obtain a b where q: "q = [:a, b:]"
    by (metis One_nat_def degree_pCons_eq_if nat.distinct(1) nat.inject
  pCons_cases)
  have "unit_factor q = 1"
    using assms by auto
  thus ?thesis
    using that[of "-a"] q <degree q = 1>
    by (auto simp: unit_factor_poly_def one_pCons dvd_field_iff is_unit_unit_factor
  split: if_splits)
qed

lemma prime_factors_alg_closed_poly_bij_betw:
  assumes "p ≠ (0 :: 'a :: {alg_closed_field, field_gcd} poly)"
  shows "bij_betw (λx. [:-x, 1:]) {x. poly p x = 0} (prime_factors p)"
proof (rule bij_betwI[of _ _ _ "λq. -poly q 0"], goal_cases)
  case 1
  have [simp]: "p div [:1:] = p" for p :: "'a poly"
    by (simp add: pCons_one)
  show ?case using assms
    by (auto simp: in_prime_factors_iff dvd_iff_poly_eq_0 prime_def

```

```

      prime_elem_linear_field_poly normalize_poly_def one_pCons)
qed (auto simp: in_prime_factors_iff elim!: prime_poly_alg_closedE dvdE)

lemma alg_closed_imp_factorization':
  assumes "p ≠ 0 :: 'a :: alg_closed_field poly"
  shows "p = smult (lead_coeff p) (∏ x | poly p x = 0. [:-x, 1:] ^ order
x p)"
proof -
  obtain A where A: "size A = degree p" "p = smult (lead_coeff p) (∏ x∈#A.
[:- x, 1:])"
  using alg_closed_imp_factorization[OF assms] by blast
  have "set_mset A = {x. poly p x = 0}" using assms
  by (subst A(2)) (auto simp flip: poly_hom.prod_mset_image simp: image_image)

  note A(2)
  also have "(∏ x∈#A. [:- x, 1:]) =
(∏ x∈(λx. [:- x, 1:]) ` set_mset A. x ^ count {#[:- x,
1:]. x ∈# A#} x)"
  by (subst prod_mset_multiplicity) simp_all
  also have "set_mset A = {x. poly p x = 0}" using assms
  by (subst A(2)) (auto simp flip: poly_hom.prod_mset_image simp: image_image)
  also have "(∏ x∈(λx. [:- x, 1:]) ` {x. poly p x = 0}. x ^ count {#[:-
x, 1:]. x ∈# A#} x) =
(∏ x | poly p x = 0. [:- x, 1:] ^ count {#[:- x, 1:]. x ∈#
A#} [:- x, 1:])"
  by (subst prod.reindex) (auto intro: inj_onI)
  also have "(λx. count {#[:- x, 1:]. x ∈# A#} [:- x, 1:]) = count A"
  by (subst count_image_mset_inj) (auto intro!: inj_onI)
  also have "count A = (λx. order x p)"
proof
  fix x :: 'a
  have "order x p = order x (∏ x∈#A. [:- x, 1:])"
  using assms by (subst A(2)) (auto simp: order_smult order_prod_mset)
  also have "... = (∑ y∈#A. order x [:-y, 1:])"
  by (subst order_prod_mset) (auto simp: multiset.map_comp o_def)
  also have "image_mset (λy. order x [:-y, 1:]) A = image_mset (λy.
if y = x then 1 else 0) A"
  using order_power_n_n[of y 1 for y :: 'a]
  by (intro image_mset_cong) (auto simp: order_0I)
  also have "... = replicate_mset (count A x) 1 + replicate_mset (size
A - count A x) 0"
  by (induction A) (auto simp: add_ac Suc_diff_le count_le_size)
  also have "sum_mset ... = count A x"
  by simp
  finally show "count A x = order x p" ..
qed
finally show ?thesis .
qed

```

1.4 Complex polynomials and conjugation

```

lemma complex_poly_real_coeffsE:
  assumes "set (coeffs p)  $\subseteq$   $\mathbb{R}$ "
  obtains p' where "p = map_poly complex_of_real p'"
proof (rule that)
  have "coeff p n  $\in$   $\mathbb{R}$ " for n
    using assms by (metis Reals_0 coeff_in_coeffs in_mono le_degree zero_poly.rep_eq)
  thus "p = map_poly complex_of_real (map_poly Re p)"
    by (subst map_poly_map_poly) (auto simp: poly_eq_iff o_def coeff_map_poly)
qed

```

```

lemma order_map_poly_cnj:
  assumes "p  $\neq$  0"
  shows "order x (map_poly cnj p) = order (cnj x) p"
proof -
  have "order x (map_poly cnj p)  $\leq$  order (cnj x) p" if p: "p  $\neq$  0" for
  p :: "complex poly" and x
  proof (rule order_max)
    interpret map_poly_idom_hom cnj
      by standard auto
    interpret field_hom cnj
      by standard auto
    have "[: -x, 1:] ^ order x (map_poly cnj p) dvd map_poly cnj p"
      using order[of "map_poly cnj p" x] p by simp
    also have "[: -x, 1:] ^ order x (map_poly cnj p) =
      map_poly cnj ([: -cnj x, 1:] ^ order x (map_poly cnj p))"
      by (simp add: hom_power)
    finally show "[: -cnj x, 1:] ^ order x (map_poly cnj p) dvd p"
      by (rule dvd_map_poly_hom_imp_dvd)
  qed fact+
  from this[of p x] and this[of "map_poly cnj p" "cnj x"] and assms show
  ?thesis
  by (simp add: map_poly_map_poly o_def)
qed

```

1.5 n-ary product rule for the derivative

```

lemma has_field_derivative_prod_mset [derivative_intros]:
  assumes " $\bigwedge x. x \in \# A \implies (f x \text{ has\_field\_derivative } f' x) \text{ (at } z)$ "
  shows " $((\lambda u. \prod_{x \in \# A}. f x u) \text{ has\_field\_derivative } (\sum_{x \in \# A}. f' x * (\prod_{y \in \# A - \{x\}}. f y z))) \text{ (at } z)$ "
  using assms
proof (induction A)
  case (add x A)
  note [derivative_intros] = add
  note [cong] = image_mset_cong_simp
  show ?case
    by (auto simp: field_simps multiset.map_comp o_def intro!: derivative_eq_intros)
qed auto

```

```

lemma has_field_derivative_prod [derivative_intros]:
  assumes " $\bigwedge x. x \in A \implies (f\ x\ \text{has\_field\_derivative}\ f'\ x)\ (\text{at}\ z)$ "
  shows " $((\lambda u. \prod_{x \in A}. f\ x\ u)\ \text{has\_field\_derivative}\ (\sum_{x \in A}. f'\ x * (\prod_{y \in A - \{x\}}. f\ y\ z)))\ (\text{at}\ z)$ "
  using assms
proof (cases "finite A")
  case [simp, intro]: True
  have " $((\lambda u. \prod_{x \in A}. f\ x\ u)\ \text{has\_field\_derivative}\ (\sum_{x \in A}. f'\ x * (\prod_{y \in \#mset\_set\ A - \{x\}}. f\ y\ z)))\ (\text{at}\ z)$ "
    using has_field_derivative_prod_mset[of "mset_set A" f f' z] assms
    by (simp add: prod_unfold_prod_mset sum_unfold_sum_mset)
  also have " $(\sum_{x \in A}. f'\ x * (\prod_{y \in \#mset\_set\ A - \{x\}}. f\ y\ z)) = (\sum_{x \in A}. f'\ x * (\prod_{y \in \#mset\_set\ (A - \{x\}}. f\ y\ z))$ "
    by (intro sum.cong) (auto simp: mset_set_Diff)
  finally show ?thesis
    by (simp add: prod_unfold_prod_mset)
qed auto

```

```

lemma has_field_derivative_prod_mset':
  assumes " $\bigwedge x. x \in\# A \implies f\ x\ z \neq 0$ "
  assumes " $\bigwedge x. x \in\# A \implies (f\ x\ \text{has\_field\_derivative}\ f'\ x)\ (\text{at}\ z)$ "
  defines "P  $\equiv (\lambda A\ u. \prod_{x \in\# A}. f\ x\ u)$ "
  shows " $(P\ A\ \text{has\_field\_derivative}\ (P\ A\ z * (\sum_{x \in\# A}. f'\ x / f\ x\ z)))\ (\text{at}\ z)$ "
  using assms
  by (auto intro!: derivative_eq_intros cong: image_mset_cong_simp
      simp: sum_distrib_right mult_ac prod_mset_diff image_mset_Diff
      multiset.map_comp o_def)

```

```

lemma has_field_derivative_prod':
  assumes " $\bigwedge x. x \in A \implies f\ x\ z \neq 0$ "
  assumes " $\bigwedge x. x \in A \implies (f\ x\ \text{has\_field\_derivative}\ f'\ x)\ (\text{at}\ z)$ "
  defines "P  $\equiv (\lambda A\ u. \prod_{x \in A}. f\ x\ u)$ "
  shows " $(P\ A\ \text{has\_field\_derivative}\ (P\ A\ z * (\sum_{x \in A}. f'\ x / f\ x\ z)))\ (\text{at}\ z)$ "
proof (cases "finite A")
  case True
  show ?thesis using assms True
    by (auto intro!: derivative_eq_intros
        simp: prod_diff1 sum_distrib_left sum_distrib_right mult_ac)
qed (auto simp: P_def)

```

1.6 Facts about complex numbers

```

lemma Re_sum_mset: "Re (sum_mset X) =  $(\sum_{x \in\# X}. \text{Re}\ x)$ "
  by (induction X) auto

```

```

lemma Im_sum_mset: "Im (sum_mset X) =  $(\sum_{x \in\# X}. \text{Im}\ x)$ "

```



```

    by (induction X) auto

lemma Re_sum_mset': "Re ( $\sum x \in \#X. f x$ ) = ( $\sum x \in \#X. \text{Re } (f x)$ )"
  by (induction X) auto

lemma Im_sum_mset': "Im ( $\sum x \in \#X. f x$ ) = ( $\sum x \in \#X. \text{Im } (f x)$ )"
  by (induction X) auto

lemma inverse_complex_altdef: "inverse z = cnj z / norm z ^ 2"
  by (metis complex_div_cnj inverse_eq_divide mult_1)

end

theory Polynomial_Crit_Geometry
imports
  "HOL-Computational_Algebra.Computational_Algebra"
  "HOL-Analysis.Analysis"
  Polynomial_Crit_Geometry_Library
begin

```

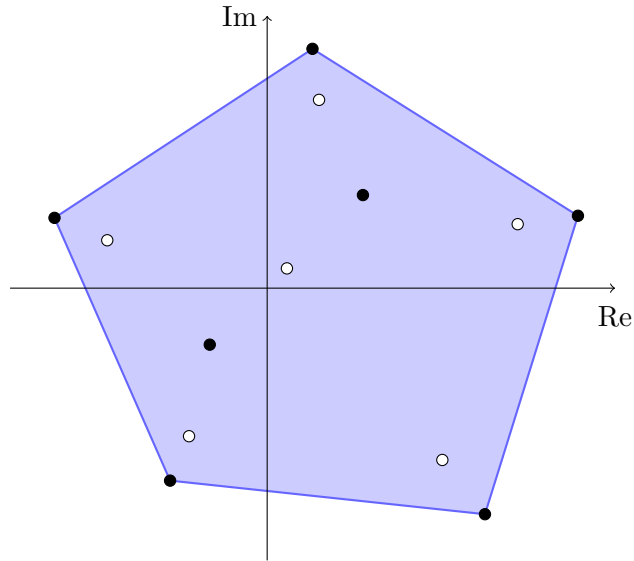


Figure 1: Example for the Gauß–Lucas Theorem: The roots (●) and critical points (○) of $x^7 - 2x^6 + x^5 + x^4 - (1 + i)x^3 - 15ix^2 - 4(1 - i)x - 7$. The critical points all lie inside the convex hull of the roots (□).

2 The Gauß–Lucas Theorem

The following result is known as the *Gauß–Lucas Theorem*: The critical points of a non-constant complex polynomial lie inside the convex hull of its roots.

The proof is relatively straightforward by writing the polynomial in the form

$$p(x) = \prod_{i=1}^n (x - x_i)^{a_i} ,$$

from which we get the derivative

$$p'(x) = p(x) \cdot \sum_{i=1}^n \frac{a_i}{x - x_i} .$$

With some more calculations, one can then see that every root x of p' can be written as

$$x = \sum_{i=1}^n \frac{u_i}{U} \cdot x_i$$

where $u_i = \frac{a_i}{|x - x_i|^2}$ and $U = \sum_{i=1}^n u_i$.

theorem `pderiv_roots_in_convex_hull`:

```

fixes p :: "complex poly"
assumes "degree p ≠ 0"
shows "{z. poly (pderiv p) z = 0} ⊆ convex hull {z. poly p z = 0}"
proof safe
  fix z :: complex
  assume "poly (pderiv p) z = 0"
  show "z ∈ convex hull {z. poly p z = 0}"
  proof (cases "poly p z = 0")
    case True
      thus ?thesis by (simp add: hull_inc)
    next
      case False
        hence [simp]: "p ≠ 0" by auto
        define α where "α = lead_coeff p"
        have p_eq: "p = smult α (∏ z | poly p z = 0. [:- z, 1:] ^ order z
p)"
          unfolding α_def by (rule alg_closed_imp_factorization') fact
          have poly_p: "poly p = (λw. α * (∏ z | poly p z = 0. (w - z) ^ order
z p))"
            by (subst p_eq) (simp add: poly_prod fun_eq_iff)

          define S where "S = (∑ w | poly p w = 0. of_nat (order w p) / (z
- w))"
          define u :: "complex ⇒ real" where "u = (λw. of_nat (order w p) /
norm (z - w) ^ 2)"
          define U where "U = (∑ w | poly p w = 0. u w)"
          have u_pos: "u w > 0" if "poly p w = 0" for w
            using that False by (auto simp: u_def order_pos_iff intro!: divide_pos_pos)
          hence "U > 0" unfolding U_def
            using assms fundamental_theorem_of_algebra[of p] False
            by (intro sum_pos poly_roots_finite) (auto simp: constant_degree)

          note [derivative_intros del] = has_field_derivative_prod
          note [derivative_intros] = has_field_derivative_prod'
          have "(poly p has_field_derivative poly p z *
(∑ w | poly p w = 0. of_nat (order w p) *
(z - w) ^ (order w p - 1) / (z - w) ^ order w p) ) (at
z)"
            (is "(_ has_field_derivative _ * ?S') _") using False
            by (subst (1 2) poly_p)
            (auto intro!: derivative_eq_intros simp: order_pos_iff mult_ac
power_diff S_def)
          also have "?S' = S" unfolding S_def
          proof (intro sum.cong refl, goal_cases)
            case (1 w)
              with False have "w ≠ z" and "order w p > 0"
                by (auto simp: order_pos_iff)
              thus ?case by (simp add: power_diff)
          qed

```

```

finally have "(poly p has_field_derivative poly p z * S) (at z)" .
hence "poly (pderiv p) z = poly p z * S"
  by (rule sym[OF DERIV_unique]) (auto intro: poly_DERIV)
with <poly (pderiv p) z = 0> and <poly p z ≠ 0> have "S = 0" by
simp

also have "S = (∑ w | poly p w = 0. of_nat (order w p) * cnj z / norm
(z - w) ^ 2 -
of_nat (order w p) * cnj w / norm
(z - w) ^ 2)"
  unfolding S_def by (intro sum.cong refl, subst complex_div_cnj)
  (auto simp: diff_divide_distrib ring_distrib)
also have "... = cnj z * (∑ w | poly p w = 0. u w) - (∑ w | poly p
w = 0. u w * cnj w)"
  by (simp add: sum_subtractf sum_distrib_left mult_ac u_def)
finally have "cnj z * (∑ w | poly p w = 0. of_real (u w)) =
(∑ w | poly p w = 0. of_real (u w) * cnj w)" by simp
from arg_cong[OF this, of cnj]
have "z * of_real U = (∑ w | poly p w = 0. of_real (u w) * w)"
  unfolding complex_cnj_mult by (simp add: U_def)
hence "z = (∑ w | poly p w = 0. of_real (u w) * w) / of_real U"
  using <U > 0> by (simp add: divide_simps)
also have "... = (∑ w | poly p w = 0. (u w / U) *R w)"
  by (subst sum_divide_distrib) (auto simp: scaleR_conv_of_real)
finally have z_eq: "z = (∑ w | poly p w = 0. (u w / U) *R w)" .

show "z ∈ convex hull {z. poly p z = 0}"
proof (subst z_eq, rule convex_sum)
  have "(∑ i∈{w. poly p w = 0}. u i / U) = U / U"
    by (subst (2) U_def) (simp add: sum_divide_distrib)
  also have "... = 1" using <U > 0> by simp
  finally show "(∑ i∈{w. poly p w = 0}. u i / U) = 1" .
qed (insert <U > 0> u_pos,
  auto simp: hull_inc intro!: divide_nonneg_pos less_imp_le poly_roots_finite)
qed
qed

```

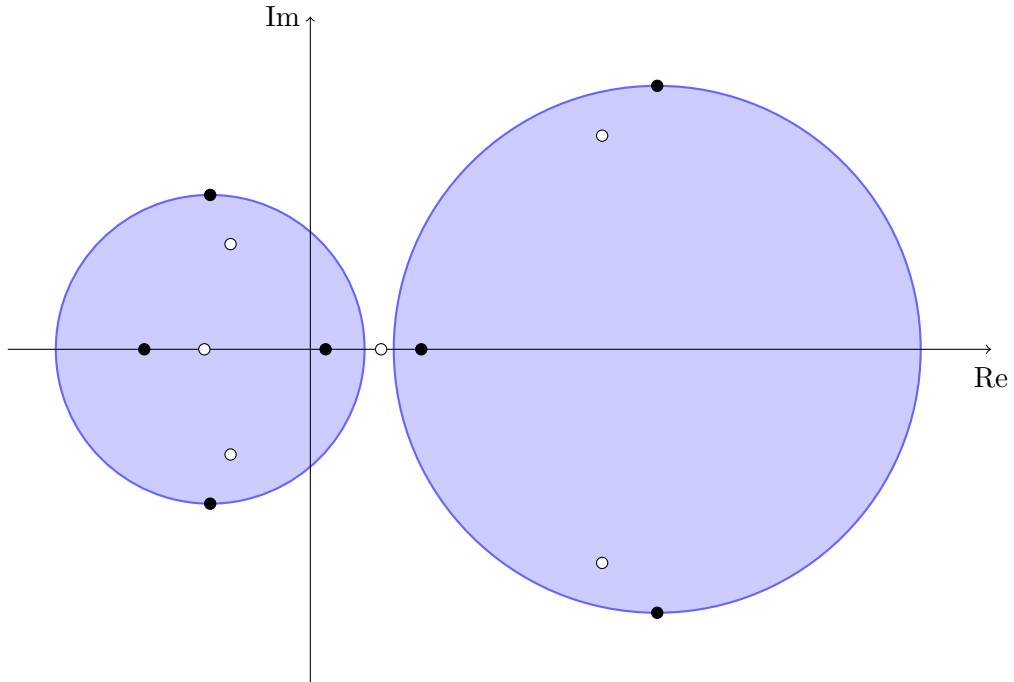


Figure 2: Example for Jensen's Theorem: The roots (\bullet) and critical points (\circ) of the polynomial $x^7 - 3x^6 + 2x^5 + 8x^4 + 10x^3 - 10x + 1$. It can be seen that all the non-real critical points lie inside a Jensen disc (\circ), whereas there can be real critical points that do *not* lie inside a Jensen disc.

3 Jensen's Theorem

For each root w of a real polynomial p , the Jensen disc of w is the smallest disc containing both w and \bar{w} , i.e. the disc with centre $\text{Re}(w)$ and radius $|\text{Im}(w)|$.

We now show that if p is a real polynomial, every non-real critical point of p lies inside a Jensen disc of one of its non-real roots.

definition `jensen_disc` :: "complex \Rightarrow complex set" where
`"jensen_disc w = cball (of_real (Re w)) |Im w|"`

theorem `pderiv_root_in_jensen_disc`:
`fixes p :: "complex poly"`
`assumes "set (coeffs p) \subseteq \mathbb{R} " and "degree p \neq 0"`
`assumes "poly (pderiv p) z = 0" and "z \notin \mathbb{R} "`
`shows " $\exists w. w \notin \mathbb{R} \wedge \text{poly } p w = 0 \wedge z \in \text{jensen_disc } w$ "`
proof (`rule ccontr`)
`have real_coefs: "coeff p n \in \mathbb{R} " for n`

```

    using assms(1) by (metis Reals_0 coeff_0 coeff_in_coeffs le_degree
subsetD)
  define d where "d = (λx. dist z (Re x) ^ 2 - Im x ^ 2)"

  assume *: "¬(∃w. w ∉ ℝ ∧ poly p w = 0 ∧ z ∈ jensen_disc w)"
  have d_pos: "d w > 0" if "poly p w = 0" "w ∉ ℝ" for w
  proof -
    have "dist z (Re w) > |Im w|"
      using * that unfolding d_def jensen_disc_def by (auto simp: dist_commute)
    hence "dist z (Re w) ^ 2 > |Im w| ^ 2"
      by (intro power_strict_mono) auto
    thus ?thesis
      by (simp add: d_def)
  qed

  have "poly p z ≠ 0"
    using d_pos[of z] assms by (auto simp: d_def dist_norm cmod_power2)
  hence [simp]: "p ≠ 0" by auto
  define α where "α = lead_coeff p"
  have [simp]: "α ≠ 0"
    using assms(4) by (auto simp: α_def)
  obtain A where p_eq: "p = smult α (∏ x∈#A. [:-x, 1:])"
    unfolding α_def using alg_closed_imp_factorization[of p] by auto
  have poly_p: "poly p = (λw. α * (∏ z∈#A. w - z))"
    by (subst p_eq) (simp add: poly_prod_mset fun_eq_iff)
  have [simp]: "poly p z = 0 ⟷ z ∈# A" for z
    by (auto simp: poly_p α_def)

  define Apos where "Apos = filter_mset (λw. Im w > 0) A"
  define Aneg where "Aneg = filter_mset (λw. Im w < 0) A"
  define A0 where "A0 = filter_mset (λw. Im w = 0) A"
  have "A = Apos + Aneg + A0"
    unfolding Apos_def Aneg_def A0_def by (induction A) auto

  have count_A: "count A w = order w p" for w
  proof -
    have "0 ∉# {#[:- x, 1:]. x ∈# A#}"
      by auto
    hence "order w p = (∑ x∈#A. order w [:- x, 1:])"
      by (simp add: p_eq order_smult order_prod_mset multiset.map_comp
o_def)
    also have "... = (∑ x∈#A. if w = x then 1 else 0)"
      by (simp add: order_linear_factor)
    also have "... = count A w"
      by (induction A) auto
    finally show ?thesis ..
  qed

  have "Aneg = image_mset cnj Apos"

```

```

proof (rule multiset_eqI)
  fix x :: complex
  have "order (cnj x) (map_poly cnj p) = order x p"
    by (subst order_map_poly_cnj) auto
  also have "map_poly cnj p = p"
    using assms(1) by (metis Reals_cnj_iff map_poly_idI' subsetD)
  finally have [simp]: "order (cnj x) p = order x p" .

  have "count (image_mset cnj Apos) (cnj (cnj x)) = count Apos (cnj
x)"
    by (subst count_image_mset_inj) (auto simp: inj_on_def)
  also have "... = count Aneg x"
    by (simp add: Apos_def Aneg_def count_A)
  finally show "count Aneg x = count (image_mset cnj Apos) x"
    by simp
qed

have [simp]: "cnj x ∈# A ⟷ x ∈# A" for x
proof -
  have "cnj x ∈# A ⟷ poly p (cnj x) = 0"
    by simp
  also have "poly p (cnj x) = cnj (poly (map_poly cnj p) x)"
    by simp
  also have "map_poly cnj p = p"
    using real_coefs by (intro poly_eqI) (auto simp: coeff_map_poly
Reals_cnj_iff)
  finally show ?thesis
    by simp
qed

define N where "N = (λx. norm ((z - x) * (z - cnj x)))"
have N_pos: "N x > 0" if "x ∈# A" for x
  using that <poly p z ≠ 0> by (auto simp: N_def)
have N_nonneg: "N x ≥ 0" and [simp]: "N x ≠ 0" if "x ∈# A" for x
  using N_pos[OF that] by simp_all

We show that  $(\sum_{x \in \#A}. 1 / (z - x)) = 0$  (which is relatively obvious) and
then that the imaginary part of this sum is positive, which is a contradiction.

define S where "S =  $(\sum_{x \in \#A}. 1 / (z - x))$ "
note [derivative_intros del] = has_field_derivative_prod_mset
note [derivative_intros] = has_field_derivative_prod_mset'
have "(poly p has_field_derivative poly p z * S) (at z)"
  using <poly p z ≠ 0> unfolding S_def
  by (subst (1 2) poly_p)
  (auto intro!: derivative_eq_intros simp: order_pos_iff mult_ac
power_diff multiset.map_comp o_def)
hence "poly (pderiv p) z = poly p z * S"
  by (rule sym[OF DERIV_unique]) (auto intro: poly_DERIV)
with <poly (pderiv p) z = 0> and <poly p z ≠ 0> have "S = 0" by simp

```

For determining $\text{Im } S$, we decompose the sum into real roots and pairs of conjugate and merge the sum of each pair of conjugate roots.

```

have "Im S = (∑ x∈#Apos. Im (1 / (z - x))) + (∑ x∈#Aneg. Im (1 / (z - x))) + (∑ x∈#A0. Im (1 / (z - x)))"
  by (simp add: S_def ‹A = Apos + Aneg + A0› Im_sum_mset')
also have "Aneg = image_mset cnj Apos"
  by fact
also have "(∑ x∈#... Im (1 / (z - x))) = (∑ x∈#Apos. Im (1 / (z - cnj x)))"
  by (simp add: multiset.map_comp o_def)
also have "(∑ x∈#Apos. Im (1 / (z - x))) + (∑ x∈#Apos. Im (1 / (z - cnj x))) =
  (∑ x∈#Apos. Im (1 / (z - x) + 1 / (z - cnj x)))"
  by (subst sum_mset.distrib [symmetric]) simp_all
also have "image_mset (λx. Im (1 / (z - x) + 1 / (z - cnj x))) Apos
=
  image_mset (λx. - 2 * Im z * d x / N x ^ 2) Apos"
proof (intro image_mset_cong, goal_cases)
  case (1 x)
  have "1 / (z - x) + 1 / (z - cnj x) = (2 * z - (x + cnj x)) * inverse
((z - x) * (z - cnj x))"
    using ‹poly p z ≠ 0› 1
    by (auto simp: divide_simps Apos_def complex_is_Real_iff simp flip:
Reals_cnj_iff)
  also have "x + cnj x = 2 * Re x"
    by (subst complex_add_cnj) auto
  also have "inverse ((z - x) * (z - cnj x)) = (cnj z - cnj x) * (cnj
z - x) / N x ^ 2"
    by (subst inverse_complex_altdef) (simp_all add: N_def)
  also have "Im ((2 * z - complex_of_real (2 * Re x)) * ((cnj z - cnj
x) * (cnj z - x) / N x ^ 2)) =
    (-2 * Im z * (Im z ^ 2 - Im x ^ 2 + (Re x - Re z) ^ 2))
/ N x ^ 2"
    by (simp add: algebra_simps power2_eq_square)
  also have "Im z ^ 2 - Im x ^ 2 + (Re x - Re z) ^ 2 = d x"
    unfolding dist_norm cmod_power2 d_def by (simp add: power2_eq_square
algebra_simps)
  finally show ?case .
qed
also have "sum_mset ... = -Im z * (∑ x∈#Apos. 2 * d x / N x ^ 2)"
  by (subst sum_mset_distrib_left) (simp_all add: multiset.map_comp
o_def mult_ac)
also have "image_mset (λx. Im (1 / (z - x))) A0 = image_mset (λx. -Im
z / N x) A0"
proof (intro image_mset_cong, goal_cases)
  case (1 x)
  have [simp]: "Im x = 0"
    using 1 by (auto simp: A0_def)
  have [simp]: "cnj x = x"

```



```

    by (auto simp: complex_eq_iff)
    show "Im (1 / (z - x)) = -Im z / N x"
    by (simp add: Im_divide N_def cmod_power2 norm_power flip: power2_eq_square)
qed
also have "sum_mset ... = -Im z * ( $\sum_{x \in \#A0} 1 / N x$ )"
  by (simp add: sum_mset_distrib_left multiset.map_comp o_def)
also have "-Im z * ( $\sum_{x \in \#Apos} 2 * d x / N x ^ 2$ ) + ... =
  -Im z * (( $\sum_{x \in \#Apos} 2 * d x / N x ^ 2$ ) + ( $\sum_{x \in \#A0} 1 / N x$ ))"
  by algebra
also have "Im S = 0"
  using <S = 0> by simp
finally have "(( $\sum_{x \in \#Apos} 2 * d x / N x ^ 2$ ) + ( $\sum_{x \in \#A0} 1 / N x$ ))
= 0"
  using <z  $\notin$   $\mathbb{R}$ > by (simp add: complex_is_Real_iff)

moreover have "(( $\sum_{x \in \#Apos} 2 * d x / N x ^ 2$ ) + ( $\sum_{x \in \#A0} 1 / N x$ )) > 0"
proof -
  have "A  $\neq$  {}"
    using <degree p  $\neq$  0> p_eq by fastforce
  hence "Apos  $\neq$  {}  $\vee$  A0  $\neq$  {}"
    using <Aneg = image_mset cnj Apos> <A = Apos + Aneg + A0> by auto
  thus ?thesis
proof
  assume "Apos  $\neq$  {}"
  hence "(( $\sum_{x \in \#Apos} 2 * d x / N x ^ 2$ ) > 0"
    by (intro sum_mset_pos)
    (auto intro!: mult_pos_pos divide_pos_pos d_pos simp: Apos_def
complex_is_Real_iff)
  thus ?thesis
    by (intro add_pos_nonneg sum_mset_nonneg) (auto intro!: N_nonneg
simp: A0_def)
next
  assume "A0  $\neq$  {}"
  hence "(( $\sum_{x \in \#A0} 1 / N x$ ) > 0"
    by (intro sum_mset_pos) (auto intro!: divide_pos_pos N_pos simp:
A0_def)
  thus ?thesis
    by (intro add_nonneg_pos sum_mset_nonneg)
    (auto intro!: N_pos less_imp_le[OF d_pos] mult_nonneg_nonneg
divide_nonneg_pos
simp: Apos_def complex_is_Real_iff)
qed
qed

ultimately show False
  by simp
qed

```

end

References

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