

The Poincaré-Bendixson Theorem

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Contents

1	Additions to HOL-Analysis	1
1.1	Unsorted Lemmas (TODO: sort!)	1
1.2	indexing euclidean space with natural numbers	9
1.3	derivatives	10
1.4	Segments	11
1.5	Open Segments	13
1.6	Syntax	13
1.7	Paths	13
2	Additions to the ODE Library	14
2.1	Comparison Principle	14
2.2	Locally Lipschitz ODEs	15
2.3	Reverse flow as Sublocale	15
2.4	Autonomous LL ODE : Existence Interval and trapping on the interval	16
2.5	Connectedness	18
2.6	Return Time and Implicit Function Theorem	19
2.7	Fixpoints	21
3	Invariance	21
3.1	Tools for proving invariance	23
4	Limit Sets	24
5	Periodic Orbits	28
6	Poincare Bendixson Theory	32
6.1	Flow to Path	32
6.2	2D Line segments	33
6.3	Bijection Real-Complex for Jordan Curve Theorem	35
6.4	Transversal Segments	37
6.5	Monotone Step Lemma	40

6.6	Straightening	41
6.7	Unique Intersection	42
6.8	Poincare Bendixson Theorems	43
7	Branch-And-Bound Arithmetic	45
8	Examples	46
8.1	Simple	46
8.2	Glycolysis	48

1 Additions to HOL-Analysis

```
theory Analysis-Misc
  imports
    Ordinary-Differential-Equations.ODE-Analysis
begin
```

1.1 Unsorted Lemmas (TODO: sort!)

```
lemma uminus-uminus-image: uminus ' uminus ' S = S
  for S::'r::ab-group-add set
  <proof>
```

```
lemma in-uminus-image-iff[simp]: x ∈ uminus ' S ↔ - x ∈ S
  for S::'r::ab-group-add set
  <proof>
```

```
lemma closed-subsegmentI:
  w + t *R (z - w) ∈ {x--y}
  if w ∈ {x -- y} z ∈ {x -- y} and t: 0 ≤ t ≤ 1
  <proof>
```

```
lemma tendsto-minus-cancel-right: ((λx. -g x) → l) F ↔ (g → -l) F
  - cf (?f → - ?y) ?F = ((λx. - ?f x) → ?y) ?F
  for g::- ⇒ 'b::topological-group-add
  <proof>
```

```
lemma tendsto-nhds-continuousI: (f → l) (nhds x) if (f → l) (at x) f x = l
  - TODO: the assumption is continuity of f at x
  <proof>
```

```
lemma inj-composeD:
  assumes inj (λx. g (t x))
  shows inj t
  <proof>
```

```
lemma compact-sequentialE:
  fixes S T::'a::first-countable-topology set
```

assumes *compact S*
assumes *infinite T*
assumes $T \subseteq S$
obtains $t::nat \Rightarrow 'a$ **and** $l::'a$
where $\bigwedge n. t\ n \in T \ \bigwedge n. t\ n \neq l\ t \longrightarrow l\ l \in S$
 ⟨*proof*⟩

lemma *infinite-countable-subsetE*:
fixes $S::'a$ *set*
assumes *infinite S*
obtains $g::nat \Rightarrow 'a$ **where** $inj\ g\ range\ g \subseteq S$
 ⟨*proof*⟩

lemma *real-quad-ge*: $2 * (an * bn) \leq an * an + bn * bn$ **for** $an\ bn::real$
 ⟨*proof*⟩

lemma *inner-quad-ge*: $2 * (a \cdot b) \leq a \cdot a + b \cdot b$
for $a\ b::'a::euclidean-space$ — generalize?
 ⟨*proof*⟩

lemma *inner-quad-gt*: $2 * (a \cdot b) < a \cdot a + b \cdot b$
if $a \neq b$
for $a\ b::'a::euclidean-space$ — generalize?
 ⟨*proof*⟩

lemma *closed-segment-line-hyperplanes*:
 $\{a \dashv\dashv b\} = range\ (\lambda u. a + u *_{\mathbb{R}} (b - a)) \cap \{x. a \cdot (b - a) \leq x \cdot (b - a) \wedge x$
 $\cdot (b - a) \leq b \cdot (b - a)\}$
if $a \neq b$
for $a\ b::'a::euclidean-space$
 ⟨*proof*⟩

lemma *open-segment-line-hyperplanes*:
 $\{a <-\!-\!< b\} = range\ (\lambda u. a + u *_{\mathbb{R}} (b - a)) \cap \{x. a \cdot (b - a) < x \cdot (b - a)$
 $\wedge x \cdot (b - a) < b \cdot (b - a)\}$
if $a \neq b$
for $a\ b::'a::euclidean-space$
 ⟨*proof*⟩

lemma *at-within-interior*: *NO-MATCH UNIV S* $\Longrightarrow x \in interior\ S \Longrightarrow at\ x\ within$
 $S = at\ x$
 ⟨*proof*⟩

lemma *tendsto-at-topI*:
 $(f \longrightarrow l)$ *at-top* **if** $\bigwedge e. 0 < e \Longrightarrow \exists x_0. \forall x \geq x_0. dist\ (f\ x)\ l < e$
for $f::'a::linorder-topology \Rightarrow 'b::metric-space$
 ⟨*proof*⟩

lemma *tendsto-at-topE*:

fixes $f::'a::\text{linorder-topology} \Rightarrow 'b::\text{metric-space}$
assumes $(f \longrightarrow l) \text{ at-top}$
assumes $e > 0$
obtains $x0$ **where** $\bigwedge x. x \geq x0 \implies \text{dist } (f x) l < e$
 $\langle \text{proof} \rangle$
lemma *tendsto-at-top-iff*: $(f \longrightarrow l) \text{ at-top} \longleftrightarrow (\forall e>0. \exists x0. \forall x \geq x0. \text{dist } (f x) l < e)$
for $f::'a::\text{linorder-topology} \Rightarrow 'b::\text{metric-space}$
 $\langle \text{proof} \rangle$

lemma *tendsto-at-top-eq-left*:
fixes $f g::'a::\text{linorder-topology} \Rightarrow 'b::\text{metric-space}$
assumes $(f \longrightarrow l) \text{ at-top}$
assumes $\bigwedge x. x \geq x0 \implies f x = g x$
shows $(g \longrightarrow l) \text{ at-top}$
 $\langle \text{proof} \rangle$

lemma *lim-divide-n*: $(\lambda x. e / \text{real } x) \longrightarrow 0$
 $\langle \text{proof} \rangle$

definition *at-top-within* :: $('a::\text{order}) \text{ set} \Rightarrow 'a \text{ filter}$
where *at-top-within* $s = (\text{INF } k \in s. \text{principal } (\{k \dots\} \cap s))$

lemma *at-top-within-at-top[simp]*:
shows *at-top-within UNIV* = *at-top*
 $\langle \text{proof} \rangle$

lemma *at-top-within-empty[simp]*:
shows *at-top-within* $\{\}$ = *top*
 $\langle \text{proof} \rangle$

definition *nhds-set* $X = (\text{INF } S \in \{S. \text{open } S \wedge X \subseteq S\}. \text{principal } S)$

lemma *eventually-nhds-set*:
 $(\forall_F x \text{ in nhds-set } X. P x) \longleftrightarrow (\exists S. \text{open } S \wedge X \subseteq S \wedge (\forall x \in S. P x))$
 $\langle \text{proof} \rangle$

term *filterlim* f (*nhds-set* (*frontier* X)) F — f tends to the boundary of X ?

somewhat inspired by $?l \text{ islimpt range } ?f \implies \exists r. \text{strict-mono } r \wedge (?f \circ r) \longrightarrow ?l$ and its dependencies. The class constraints seem somewhat arbitrary, perhaps this can be generalized in some way.

lemma *limpt-closed-imp-exploding-subsequence*:— TODO: improve name?!
fixes $f::'a::\{\text{heine-borel, real-normed-vector}\} \Rightarrow 'b::\{\text{first-countable-topology, t2-space}\}$
assumes *cont*[*THEN* *continuous-on-compose2*, *continuous-intros*]: *continuous-on* $T f$
assumes *closed*: *closed* T
assumes *bound*: $\bigwedge t. t \in T \implies f t \neq l$
assumes *limpt*: $l \text{ islimpt } (f ' T)$

obtains s where

$$(f \circ s) \longrightarrow l$$

$$\bigwedge i. s \ i \in T$$

$$\bigwedge C. \text{ compact } C \implies C \subseteq T \implies \forall_F i \text{ in sequentially. } s \ i \notin C$$

<proof>

lemma *Inf-islimpt: bdd-below $S \implies \text{Inf } S \notin S \implies S \neq \{\} \implies \text{Inf } S \text{ islimpt } S$ for $S::\text{real set}$*

<proof>

context *linorder*

begin

HOL-analysis doesn't seem to have these, maybe they were never needed. Some variants are around $\{?a..?b\} \cap \{?c..?d\} = \{\max ?a ?c.. \min ?b ?d\}$, but with old-style naming conventions. Change to the "modern" L. convention there?

lemma *Int-Ico[simp]:*

$$\text{shows } \{a..\} \cap \{b..\} = \{\max a b ..\}$$

<proof>

lemma *Int-Ici-Ico[simp]:*

$$\text{shows } \{a..\} \cap \{b..<c\} = \{\max a b ..<c\}$$

<proof>

lemma *Int-Ico-Ici[simp]:*

$$\text{shows } \{a..<c\} \cap \{b..\} = \{\max a b ..<c\}$$

<proof>

lemma *subset-Ico-iff[simp]:*

$$\{a..<b\} \subseteq \{c..<b\} \longleftrightarrow b \leq a \vee c \leq a$$

<proof>

lemma *Ico-subset-Ioo-iff[simp]:*

$$\{a..<b\} \subseteq \{c<..<b\} \longleftrightarrow b \leq a \vee c < a$$

<proof>

lemma *Icc-Un-Ici[simp]:*

$$\text{shows } \{a..b\} \cup \{b..\} = \{\min a b..\}$$

<proof>

end

lemma *at-top-within-at-top-unbounded-right:*

fixes *a::'a::linorder*

shows *at-top-within* $\{a..\} = \text{at-top}$

<proof>

lemma *at-top-within-at-top-unbounded-rightI:*

fixes $a::'a::\text{linorder}$
assumes $\{a..\} \subseteq s$
shows $\text{at-top-within } s = \text{at-top}$
 $\langle \text{proof} \rangle$

lemma $\text{at-top-within-at-top-bounded-right}$:
fixes $a b::'a::\{\text{dense-order},\text{linorder-topology}\}$
assumes $a < b$
shows $\text{at-top-within } \{a..<b\} = \text{at-left } b$
 $\langle \text{proof} \rangle$

lemma $\text{at-top-within-at-top-bounded-right}'$:
fixes $a b::'a::\{\text{dense-order},\text{linorder-topology}\}$
assumes $a < b$
shows $\text{at-top-within } \{..<b\} = \text{at-left } b$
 $\langle \text{proof} \rangle$

lemma $\text{eventually-at-top-within-linorder}$:
assumes $sn:s \neq \{\}$
shows $\text{eventually } P (\text{at-top-within } s) \longleftrightarrow (\exists x0::'a::\{\text{linorder-topology}\} \in s. \forall x \geq x0. x \in s \longrightarrow P x)$
 $\langle \text{proof} \rangle$

lemma $\text{tendsto-at-top-withinI}$:
fixes $f::'a::\text{linorder-topology} \Rightarrow 'b::\text{metric-space}$
assumes $s \neq \{\}$
assumes $\bigwedge e. 0 < e \implies \exists x0 \in s. \forall x \in \{x0..\} \cap s. \text{dist } (f x) l < e$
shows $(f \longrightarrow l) (\text{at-top-within } s)$
 $\langle \text{proof} \rangle$

lemma $\text{tendsto-at-top-withinE}$:
fixes $f::'a::\text{linorder-topology} \Rightarrow 'b::\text{metric-space}$
assumes $s \neq \{\}$
assumes $(f \longrightarrow l) (\text{at-top-within } s)$
assumes $e > 0$
obtains $x0$ **where** $x0 \in s \bigwedge x. x \in \{x0..\} \cap s \implies \text{dist } (f x) l < e$
 $\langle \text{proof} \rangle$

lemma $\text{tendsto-at-top-within-iff}$:
fixes $f::'a::\text{linorder-topology} \Rightarrow 'b::\text{metric-space}$
assumes $s \neq \{\}$
shows $(f \longrightarrow l) (\text{at-top-within } s) \longleftrightarrow (\forall e>0. \exists x0 \in s. \forall x \in \{x0..\} \cap s. \text{dist } (f x) l < e)$
 $\langle \text{proof} \rangle$

lemma $\text{filterlim-at-top-at-top-within-bounded-right}$:
fixes $a b::'a::\{\text{dense-order},\text{linorder-topology}\}$
fixes $f::'a \Rightarrow \text{real}$
assumes $a < b$

shows $\text{filterlim } f \text{ at-top (at-top-within } \{..<b\}) = (f \longrightarrow \infty) \text{ (at-left } b)$
 ⟨proof⟩

Extract a sequence (going to infinity) bounded away from l

lemma *not-tendsto-frequentlyE*:
assumes $\neg((f \longrightarrow l) F)$
obtains S **where** $\text{open } S \ l \in S \ \exists_F x \text{ in } F. f x \notin S$
 ⟨proof⟩

lemma *not-tendsto-frequently-metricE*:
assumes $\neg((f \longrightarrow l) F)$
obtains e **where** $e > 0 \ \exists_F x \text{ in } F. e \leq \text{dist } (f x) l$
 ⟨proof⟩

lemma *eventually-frequently-conj*: $\text{frequently } P \ F \implies \text{eventually } Q \ F \implies \text{frequently } (\lambda x. P x \wedge Q x) \ F$
 ⟨proof⟩

lemma *frequently-at-top*:
 $(\exists_F t \text{ in } \text{at-top}. P t) \longleftrightarrow (\forall t0. \exists t > t0. P t)$
for $P::'a::\{\text{linorder}, \text{no-top}\} \Rightarrow \text{bool}$
 ⟨proof⟩

lemma *frequently-at-topE*:
fixes $P::\text{nat} \Rightarrow 'a::\{\text{linorder}, \text{no-top}\} \Rightarrow -$
assumes $\text{freq}[\text{rule-format}]: \forall n. \exists_F a \text{ in } \text{at-top}. P n a$
obtains $s::\text{nat} \Rightarrow 'a$
where $\bigwedge i. P i (s i) \text{ strict-mono } s$
 ⟨proof⟩

lemma *frequently-at-topE'*:
fixes $P::\text{nat} \Rightarrow 'a::\{\text{linorder}, \text{no-top}\} \Rightarrow -$
assumes $\text{freq}[\text{rule-format}]: \forall n. \exists_F a \text{ in } \text{at-top}. P n a$
and $g: \text{filterlim } g \text{ at-top sequentially}$
obtains $s::\text{nat} \Rightarrow 'a$
where $\bigwedge i. P i (s i) \text{ strict-mono } s \ \bigwedge n. g n \leq s n$
 ⟨proof⟩

lemma *frequently-at-top-at-topE*:
fixes $P::\text{nat} \Rightarrow 'a::\{\text{linorder}, \text{no-top}\} \Rightarrow -$ **and** $g::\text{nat} \Rightarrow 'a$
assumes $\forall n. \exists_F a \text{ in } \text{at-top}. P n a \ \text{filterlim } g \text{ at-top sequentially}$
obtains $s::\text{nat} \Rightarrow 'a$
where $\bigwedge i. P i (s i) \text{ filterlim } s \text{ at-top sequentially}$
 ⟨proof⟩

lemma *not-tendsto-convergent-seq*:
fixes $f::\text{real} \Rightarrow 'a::\text{metric-space}$
assumes $X: \text{compact } (X::'a \text{ set})$

assumes *im*: $\bigwedge x. x \geq 0 \implies f x \in X$
assumes *nl*: $\neg ((f \longrightarrow (l::'a)) \text{ at-top})$
obtains *s k* **where**
 $k \in X \ k \neq l \ (f \circ s) \longrightarrow k \text{ strict-mono } s \ \forall n. \ s \ n \geq n$
 <proof>

lemma *harmonic-bound*:
shows $1 / 2 \ \frown (Suc \ n) < 1 / \text{real } (Suc \ n)$
 <proof>

lemma *INF-bounded-imp-convergent-seq*:
fixes *f*::*real* \Rightarrow *real*
assumes *cont*: *continuous-on* $\{a..\}$ *f*
assumes *bound*: $\bigwedge t. t \geq a \implies f \ t > l$
assumes *inf*: $(INF \ t \in \{a..\}. f \ t) = l$
obtains *s* **where**
 $(f \circ s) \longrightarrow l$
 $\bigwedge i. \ s \ i \in \{a..\}$
filterlim s at-top sequentially
 <proof>

lemma *filterlim-at-top-strict-mono*:
fixes *s* :: - \Rightarrow *'a*::*linorder*
fixes *r* :: *nat* \Rightarrow -
assumes *strict-mono s*
assumes *strict-mono r*
assumes *filterlim s at-top F*
shows *filterlim (s o r) at-top F*
 <proof>

lemma *LIMSEQ-lb*:
assumes *fl*: $s \longrightarrow (l::\text{real})$
assumes *u*: $l < u$
shows $\exists n0. \ \forall n \geq n0. \ s \ n < u$
 <proof>

lemma *filterlim-at-top-choose-lower*:
assumes *filterlim s at-top sequentially*
assumes $(f \circ s) \longrightarrow l$
obtains *t* **where**
filterlim t at-top sequentially
 $(f \circ t) \longrightarrow l$
 $\forall n. \ t \ n \geq (b::\text{real})$
 <proof>

lemma *frequently-at-top-realE*:
fixes *P*::*nat* \Rightarrow *real* \Rightarrow *bool*

assumes $\forall n. \exists_F t \text{ in } at\text{-top}. P n t$
obtains $s::nat \Rightarrow real$
where $\bigwedge i. P i (s i) \text{ filterlim } s \text{ at-top } at\text{-top}$
 $\langle proof \rangle$

lemma *approachable-sequenceE*:
fixes $f::real \Rightarrow 'a::metric\text{-space}$
assumes $\bigwedge t e. 0 \leq t \implies 0 < e \implies \exists tt \geq t. dist (f tt) p < e$
obtains $s \text{ where } \text{filterlim } s \text{ at-top sequentially } (f \circ s) \longrightarrow p$
 $\langle proof \rangle$

lemma *mono-inc-bdd-above-has-limit-at-topI*:
fixes $f::real \Rightarrow real$
assumes *mono* f
assumes $\bigwedge x. f x \leq u$
shows $\exists l. (f \longrightarrow l) \text{ at-top}$
 $\langle proof \rangle$

lemma *gen-mono-inc-bdd-above-has-limit-at-topI*:
fixes $f::real \Rightarrow real$
assumes $\bigwedge x y. x \geq b \implies x \leq y \implies f x \leq f y$
assumes $\bigwedge x. x \geq b \implies f x \leq u$
shows $\exists l. (f \longrightarrow l) \text{ at-top}$
 $\langle proof \rangle$

lemma *gen-mono-dec-bdd-below-has-limit-at-topI*:
fixes $f::real \Rightarrow real$
assumes $\bigwedge x y. x \geq b \implies x \leq y \implies f x \geq f y$
assumes $\bigwedge x. x \geq b \implies f x \geq u$
shows $\exists l. (f \longrightarrow l) \text{ at-top}$
 $\langle proof \rangle$

lemma *infdist-closed*:
shows *closed* $(\{z. \text{infdist } z S \geq e\})$
 $\langle proof \rangle$

lemma *LIMSEQ-norm-0-pow*:
assumes $k > 0 b > 1$
assumes $\bigwedge n::nat. norm (s n) \leq k / b^n$
shows $s \longrightarrow 0$
 $\langle proof \rangle$

lemma *filterlim-apply-filtermap*:
assumes $g: \text{filterlim } g G F$
shows $\text{filterlim } (\lambda x. m (g x)) (\text{filtermap } m G) F$
 $\langle proof \rangle$

lemma *eventually-at-right-field-le*:

eventually P (at-right x) $\longleftrightarrow (\exists b > x. \forall y > x. y \leq b \longrightarrow P y)$
for $x :: 'a :: \{\text{linordered-field, linorder-topology}\}$
 $\langle \text{proof} \rangle$

1.2 indexing euclidean space with natural numbers

definition $\text{nth-eucl} :: 'a :: \text{executable-euclidean-space} \Rightarrow \text{nat} \Rightarrow \text{real}$ **where**
 $\text{nth-eucl } x \ i = x \cdot (\text{Basis-list } ! \ i)$

— TODO: why is that and some sort of lambda-eucl nowhere available?

definition $\text{lambda-eucl} :: (\text{nat} \Rightarrow \text{real}) \Rightarrow 'a :: \text{executable-euclidean-space}$ **where**
 $\text{lambda-eucl } (f :: \text{nat} \Rightarrow \text{real}) = (\sum i < \text{DIM}('a). f \ i *_{\mathbb{R}} \text{Basis-list } ! \ i)$

lemma $\text{eucl-eq-iff}: x = y \longleftrightarrow (\forall i < \text{DIM}('a). \text{nth-eucl } x \ i = \text{nth-eucl } y \ i)$
for $x \ y :: 'a :: \text{executable-euclidean-space}$
 $\langle \text{proof} \rangle$

bundle eucl-notation **begin**
notation nth-eucl (infixl $\$_e$ 90)
end
bundle no-eucl-notation **begin**
no-notation nth-eucl (infixl $\$_e$ 90)
end

unbundle eucl-notation

lemma $\text{eucl-of-list-eucl-nth}$:
 $(\text{eucl-of-list } xs :: 'a) \ \$_e \ i = xs \ ! \ i$
if $\text{length } xs = \text{DIM}('a :: \text{executable-euclidean-space})$
 $i < \text{DIM}('a)$
 $\langle \text{proof} \rangle$

lemma $\text{eucl-of-list-inner}$:
 $(\text{eucl-of-list } xs :: 'a) \cdot \text{eucl-of-list } ys = (\sum (x,y) \leftarrow \text{zip } xs \ ys. x * y)$
if $\text{length } xs = \text{DIM}('a :: \text{executable-euclidean-space})$
 $\text{length } ys = \text{DIM}('a :: \text{executable-euclidean-space})$
 $\langle \text{proof} \rangle$

lemma $\text{self-eq-eucl-of-list}: x = \text{eucl-of-list } (\text{map } (\lambda i. x \ \$_e \ i) [0..<\text{DIM}('a)])$
for $x :: 'a :: \text{executable-euclidean-space}$
 $\langle \text{proof} \rangle$

lemma $\text{inner-nth-eucl}: x \cdot y = (\sum i < \text{DIM}('a). x \ \$_e \ i * y \ \$_e \ i)$
for $x \ y :: 'a :: \text{executable-euclidean-space}$
 $\langle \text{proof} \rangle$

lemma $\text{norm-nth-eucl}: \text{norm } x = \text{L2-set } (\lambda i. x \ \$_e \ i) \ \{..<\text{DIM}('a)\}$
for $x :: 'a :: \text{executable-euclidean-space}$
 $\langle \text{proof} \rangle$

lemma *plus-nth-eucl*: $(x + y) \$_e i = x \$_e i + y \$_e i$
and *minus-nth-eucl*: $(x - y) \$_e i = x \$_e i - y \$_e i$
and *uminus-nth-eucl*: $(-x) \$_e i = -x \$_e i$
and *scaleR-nth-eucl*: $(c *_R x) \$_e i = c *_R (x \$_e i)$
<proof>

lemma *inf-nth-eucl*: $\text{inf } x \ y \ \$_e i = \min (x \$_e i) (y \$_e i)$
if $i < \text{DIM}('a)$
for $x::'a::\text{executable-euclidean-space}$
<proof>

lemma *sup-nth-eucl*: $\text{sup } x \ y \ \$_e i = \max (x \$_e i) (y \$_e i)$
if $i < \text{DIM}('a)$
for $x::'a::\text{executable-euclidean-space}$
<proof>

lemma *le-iff-le-nth-eucl*: $x \leq y \iff (\forall i < \text{DIM}('a). (x \$_e i) \leq (y \$_e i))$
for $x::'a::\text{executable-euclidean-space}$
<proof>

lemma *eucl-less-iff-less-nth-eucl*: $\text{eucl-less } x \ y \iff (\forall i < \text{DIM}('a). (x \$_e i) < (y \$_e i))$
for $x::'a::\text{executable-euclidean-space}$
<proof>

lemma *continuous-on-nth-eucl*[*continuous-intros*]:
continuous-on $X (\lambda x. f \ x \ \$_e i)$
if *continuous-on* $X \ f$
<proof>

1.3 derivatives

lemma *eventually-at-ne*[*intro, simp*]: $\forall_F x \text{ in } \text{at } x0. x \neq x0$
<proof>

lemma *has-vector-derivative-withinD*:
fixes $f::\text{real} \Rightarrow 'b::\text{euclidean-space}$
assumes ($f \text{ has-vector-derivative } f'$) (*at* $x0$ *within* S)
shows $((\lambda x. (f \ x - f \ x0) /_R (x - x0)) \longrightarrow f')$ (*at* $x0$ *within* S)
<proof>

A *path-connected* set S entering both T and $-T$ must cross the frontier of T

lemma *path-connected-frontier*:
fixes $S :: 'a::\text{real-normed-vector set}$
assumes *path-connected* S
assumes $S \cap T \neq \{\}$
assumes $S \cap -T \neq \{\}$
obtains s **where** $s \in S \ s \in \text{frontier } T$

<proof>

lemma *path-connected-not-frontier-subset*:

fixes $S :: 'a::\text{real-normed-vector set}$

assumes *path-connected* S

assumes $S \cap T \neq \{\}$

assumes $S \cap \text{frontier } T = \{\}$

shows $S \subseteq T$

<proof>

lemma *compact-attains-bounds*:

fixes $f::'a::\text{topological-space} \Rightarrow 'b::\text{linorder-topology}$

assumes *compact*: *compact* S

assumes *ne*: $S \neq \{\}$

assumes *cont*: *continuous-on* S f

obtains $l\ u$ **where** $l \in S\ u \in S \wedge x. x \in S \implies f\ x \in \{f\ l .. f\ u\}$

<proof>

lemma *uniform-limit-const*[*uniform-limit-intros*]:

uniform-limit S $(\lambda x\ y. f\ x)$ $(\lambda \cdot. l)$ F **if** $(f \longrightarrow l)$ F

<proof>

1.4 Segments

closed-segment throws away the order that our intuition keeps

definition *line*:: $'a::\text{real-vector} \Rightarrow 'a \Rightarrow \text{real} \Rightarrow 'a$

$(\{- \text{--} -\})$

where $\{a \text{--} b\}_u = a + u *_R (b - a)$

abbreviation *line-image* $a\ b\ U \equiv (\lambda u. \{a \text{--} b\}_u) \text{ ` } U$

notation *line-image* $(\{- \text{--} -\}) \text{ ` } \cdot$

lemma *in-closed-segment-iff-line*: $x \in \{a \text{--} b\} \longleftrightarrow (\exists c \in \{0..1\}. x = \text{line } a\ b\ c)$

<proof>

lemma *in-open-segment-iff-line*: $x \in \{a <\text{--}< b\} \longleftrightarrow (\exists c \in \{0 <..< 1\}. a \neq b \wedge x = \text{line } a\ b\ c)$

<proof>

lemma *line-convex-combination1*: $(1 - u) *_R \text{line } a\ b\ i + u *_R b = \text{line } a\ b\ (i + u - i *_R u)$

<proof>

lemma *line-convex-combination2*: $(1 - u) *_R a + u *_R \text{line } a\ b\ i = \text{line } a\ b\ (i *_R u)$

<proof>

lemma *line-convex-combination12*: $(1 - u) *_R \text{line } a\ b\ i + u *_R \text{line } a\ b\ j = \text{line } a\ b\ (i + u *_R (j - i))$

<proof>

lemma *mult-less-one-less-self*: $0 < x \implies i < 1 \implies i * x < x$ **for** $i x :: \text{real}$
 ⟨proof⟩

lemma *plus-times-le-one-lemma*: $i + u - i * u \leq 1$ **if** $i \leq 1$ $u \leq 1$ **for** $i u :: \text{real}$
 ⟨proof⟩

lemma *plus-times-less-one-lemma*: $i + u - i * u < 1$ **if** $i < 1$ $u < 1$ **for** $i u :: \text{real}$
 ⟨proof⟩

lemma *line-eq-endpoint-iff[simp]*:
 $\text{line } a \ b \ i = b \longleftrightarrow (a = b \vee i = 1)$
 $a = \text{line } a \ b \ i \longleftrightarrow (a = b \vee i = 0)$
 ⟨proof⟩

lemma *line-eq-iff[simp]*: $\text{line } a \ b \ x = \text{line } a \ b \ y \longleftrightarrow (x = y \vee a = b)$
 ⟨proof⟩

lemma *line-open-segment-iff*:
 $\{\text{line } a \ b \ i < \dots < b\} = \text{line } a \ b \ \{i < \dots < 1\}$
if $i < 1$ $a \neq b$
 ⟨proof⟩

lemma *open-segment-line-iff*:
 $\{a < \dots < \text{line } a \ b \ i\} = \text{line } a \ b \ \{0 < \dots < i\}$
if $0 < i$ $a \neq b$
 ⟨proof⟩

lemma *line-closed-segment-iff*:
 $\{\text{line } a \ b \ i \dots b\} = \text{line } a \ b \ \{i \dots 1\}$
if $i \leq 1$ $a \neq b$
 ⟨proof⟩

lemma *closed-segment-line-iff*:
 $\{a \dots \text{line } a \ b \ i\} = \text{line } a \ b \ \{0 \dots i\}$
if $0 < i$ $a \neq b$
 ⟨proof⟩

lemma *closed-segment-line-line-iff*: $\{\text{line } a \ b \ i1 \dots \text{line } a \ b \ i2\} = \text{line } a \ b \ \{i1 \dots i2\}$
if $i1 \leq i2$
 ⟨proof⟩

lemma *line-line1*: $\text{line } (\text{line } a \ b \ c) \ b \ x = \text{line } a \ b \ (c + x - c * x)$
 ⟨proof⟩

lemma *line-line2*: $\text{line } a \ (\text{line } a \ b \ c) \ x = \text{line } a \ b \ (c * x)$
 ⟨proof⟩

lemma *line-in-subsegment*:

$i1 < 1 \implies i2 < 1 \implies a \neq b \implies \text{line } a \ b \ i1 \in \{\text{line } a \ b \ i2 < \!-\!-\!< b\} \longleftrightarrow i2 < i1$
 ⟨proof⟩

lemma *line-in-subsegment2*:

$0 < i2 \implies 0 < i1 \implies a \neq b \implies \text{line } a \ b \ i1 \in \{a < \!-\!-\!< \text{line } a \ b \ i2\} \longleftrightarrow i1 < i2$
 ⟨proof⟩

lemma *line-in-open-segment-iff*[simp]:

$\text{line } a \ b \ i \in \{a < \!-\!-\!< b\} \longleftrightarrow (a \neq b \wedge 0 < i \wedge i < 1)$
 ⟨proof⟩

1.5 Open Segments

lemma *open-segment-subsegment*:

assumes $x1 \in \{x0 < \!-\!-\!< x3\}$
 $x2 \in \{x1 < \!-\!-\!< x3\}$
shows $x1 \in \{x0 < \!-\!-\!< x2\}$
 ⟨proof⟩

1.6 Syntax

abbreviation *sequentially-at-top*::(nat \Rightarrow real) \Rightarrow bool

(- \longrightarrow ∞) — the is to disambiguate syntax...

where $s \longrightarrow \infty \equiv \text{filterlim } s \text{ at-top sequentially}$

abbreviation *sequentially-at-bot*::(nat \Rightarrow real) \Rightarrow bool

(- \longrightarrow $-\infty$)

where $s \longrightarrow -\infty \equiv \text{filterlim } s \text{ at-bot sequentially}$

1.7 Paths

lemma *subpath0-linepath*:

shows $\text{subpath } 0 \ u \ (\text{linepath } t \ t') = \text{linepath } t \ (t + u * (t' - t))$
 ⟨proof⟩

lemma *linepath-image0-right-open-real*:

assumes $t < (t'::\text{real})$

shows $\text{linepath } t \ t' \ ' \ \{0..<1\} = \{t..<t'\}$

⟨proof⟩

lemma *oriented-subsegment-scale*:

assumes $x1 \in \{a < \!-\!-\!< b\}$

assumes $x2 \in \{x1 < \!-\!-\!< b\}$

obtains e **where** $e > 0 \ b - a = e *_R (x2 - x1)$

⟨proof⟩

end

2 Additions to the ODE Library

theory *ODE-Misc*

imports

Ordinary-Differential-Equations.ODE-Analysis

Analysis-Misc

begin

lemma *local-lipschitz-compact-bicomposeE:*

assumes *ll: local-lipschitz T X f*

assumes *cf: $\bigwedge x. x \in X \implies \text{continuous-on } I (\lambda t. f t x)$*

assumes *cI: compact I*

assumes *I \subseteq T*

assumes *cv: continuous-on I v*

assumes *cw: continuous-on I w*

assumes *v: v ' I \subseteq X*

assumes *w: w ' I \subseteq X*

obtains *L where L > 0 $\bigwedge x. x \in I \implies \text{dist } (f x (v x)) (f x (w x)) \leq L * \text{dist } (v x) (w x)$*

<proof>

2.1 Comparison Principle

lemma *comparison-principle-le:*

fixes *f::real \Rightarrow real \Rightarrow real*

and *$\varphi \psi$::real \Rightarrow real*

assumes *ll: local-lipschitz X Y f*

assumes *cf: $\bigwedge x. x \in Y \implies \text{continuous-on } \{a..b\} (\lambda t. f t x)$*

assumes *abX: $\{a..b\} \subseteq X$*

assumes *φ' : $\bigwedge x. x \in \{a..b\} \implies (\varphi \text{ has-real-derivative } \varphi' x) (at x)$*

assumes *ψ' : $\bigwedge x. x \in \{a..b\} \implies (\psi \text{ has-real-derivative } \psi' x) (at x)$*

assumes *φ -in: $\varphi ' \{a..b\} \subseteq Y$*

assumes *ψ -in: $\psi ' \{a..b\} \subseteq Y$*

assumes *init: $\varphi a \leq \psi a$*

assumes *defect: $\bigwedge x. x \in \{a..b\} \implies \varphi' x - f x (\varphi x) \leq \psi' x - f x (\psi x)$*

shows *$\forall x \in \{a..b\}. \varphi x \leq \psi x$ (is ?th1)*

<proof>

lemma *local-lipschitz-mult:*

shows *local-lipschitz (UNIV::real set) (UNIV::real set) (*)*

<proof>

lemma *comparison-principle-le-linear:*

fixes *φ :: real \Rightarrow real*

assumes *continuous-on $\{a..b\} g$*

assumes *$(\bigwedge t. t \in \{a..b\} \implies (\varphi \text{ has-real-derivative } \varphi' t) (at t))$*

assumes *$\varphi a \leq 0$*

assumes *$(\bigwedge t. t \in \{a..b\} \implies \varphi' t \leq g t *_R \varphi t)$*

shows *$\forall t \in \{a..b\}. \varphi t \leq 0$*

<proof>

2.2 Locally Lipschitz ODEs

context *ll-on-open-it* **begin**

lemma *flow-lipschitzE*:

assumes $\{a .. b\} \subseteq \text{existence-ivl } t0 \ x$

obtains L **where** *L-lipschitz-on* $\{a .. b\}$ (*flow* $t0 \ x$)

<proof>

lemma *flow-undefined0*: $t \notin \text{existence-ivl } t0 \ x \implies \text{flow } t0 \ x \ t = 0$

<proof>

lemma *csols-undefined*: $x \notin X \implies \text{csols } t0 \ x = \{\}$

<proof>

lemmas *existence-ivl-undefined* = *existence-ivl-empty2*

end

2.3 Reverse flow as Sublocale

lemma *range-preflect-0[simp]*: $\text{range } (\text{preflect } 0) = \text{UNIV}$

<proof>

lemma *range-uminus[simp]*: $\text{range } \text{uminus} = (\text{UNIV}::'a::\text{ab-group-add set})$

<proof>

context *auto-ll-on-open* **begin**

sublocale *rev*: *auto-ll-on-open* $-f$ **rewrites** $-(-f) = f$

<proof>

lemma *existence-ivl-eq-rev0*: $\text{existence-ivl0 } y = \text{uminus } ' \text{rev.existence-ivl0 } y$ **for** y

<proof>

lemma *rev-existence-ivl-eq0*: $\text{rev.existence-ivl0 } y = \text{uminus } ' \text{existence-ivl0 } y$ **for** y

<proof>

lemma *flow-eq-rev0*: $\text{flow0 } y \ t = \text{rev.flow0 } y \ (-t)$ **for** $y \ t$

<proof>

lemma *rev-eq-flow*: $\text{rev.flow0 } y \ t = \text{flow0 } y \ (-t)$ **for** $y \ t$

<proof>

lemma *rev-flow-image-eq*: $\text{rev.flow0 } x \ ' \ S = \text{flow0 } x \ ' (\text{uminus } ' \ S)$

<proof>

lemma *flow-image-eq-rev*: $\text{flow0 } x \ ' \ S = \text{rev.flow0 } x \ ' (\text{uminus } ' \ S)$

<proof>

end

context *c1-on-open* **begin**

sublocale *rev: c1-on-open* $-f -f'$ **rewrites** $-(-f) = f$ **and** $-(-f') = f'$
<proof>

end

context *c1-on-open-euclidean* **begin**

sublocale *rev: c1-on-open-euclidean* $-f -f'$ **rewrites** $-(-f) = f$ **and** $-(-f') = f'$
<proof>

end

2.4 Autonomous LL ODE : Existence Interval and trapping on the interval

lemma *bdd-above-is-intervalI: bdd-above I*
if *is-interval* $I a \leq b$ $a \in I$ $b \notin I$ **for** $I::\text{real set}$
<proof>

lemma *bdd-below-is-intervalI: bdd-below I*
if *is-interval* $I a \leq b$ $a \notin I$ $b \in I$ **for** $I::\text{real set}$
<proof>

context *auto-ll-on-open* **begin**

lemma *open-existence-ivl0:*
assumes $x : x \in X$
shows $\exists a b. a < 0 \wedge 0 < b \wedge \{a..b\} \subseteq \text{existence-ivl0 } x$
<proof>

lemma *open-existence-ivl':*
assumes $x : x \in X$
obtains a **where** $a > 0$ $\{-a..a\} \subseteq \text{existence-ivl0 } x$
<proof>

lemma *open-existence-ivl-on-compact:*
assumes $C: C \subseteq X$ **and** *compact* C $C \neq \{\}$
obtains a **where** $a > 0 \wedge x. x \in C \implies \{-a..a\} \subseteq \text{existence-ivl0 } x$
<proof>

definition *trapped-forward* $x K \longleftrightarrow (\text{flow0 } x \text{ ' } (\text{existence-ivl0 } x \cap \{0..\}) \subseteq K)$
— TODO: use this for backwards trapped, invariant, and all assumptions

definition $trapped\text{-backward } x K \longleftrightarrow (flow0\ x \text{ ' } (existence\text{-ivl0 } x \cap \{..0\}) \subseteq K)$

definition $trapped\ x K \longleftrightarrow trapped\text{-forward } x K \wedge trapped\text{-backward } x K$

lemma $trapped\text{-iff-on-existence-ivl0}$:

$trapped\ x K \longleftrightarrow (flow0\ x \text{ ' } (existence\text{-ivl0 } x) \subseteq K)$

$\langle proof \rangle$

end

context $auto\text{-ll-on-open begin}$

lemma $infinite\text{-rev-existence-ivl0-rewrites}$:

$\{0..\} \subseteq rev.\text{existence-ivl0 } x \longleftrightarrow \{..0\} \subseteq \text{existence-ivl0 } x$

$\{..0\} \subseteq rev.\text{existence-ivl0 } x \longleftrightarrow \{0..\} \subseteq \text{existence-ivl0 } x$

$\langle proof \rangle$

lemma $trapped\text{-backward-iff-rev-trapped-forward}$:

$trapped\text{-backward } x K \longleftrightarrow rev.\text{trapped-forward } x K$

$\langle proof \rangle$

If solution is trapped in a compact set at some time on its existence interval then it is trapped forever

lemma $trapped\text{-sol-right}$:

— TODO: when building on `afp-devel` (??? outdated): <https://bitbucket.org/isa-afp/afp-devel/commits/0c3edf9248d5389197f248c723b625c419e4d3eb>

assumes $compact\ K\ K \subseteq X$

assumes $x \in X\ trapped\text{-forward } x K$

shows $\{0..\} \subseteq \text{existence-ivl0 } x$

$\langle proof \rangle$

lemma $trapped\text{-sol-right-gen}$:

assumes $compact\ K\ K \subseteq X$

assumes $t \in \text{existence-ivl0 } x\ trapped\text{-forward } (flow0\ x\ t)\ K$

shows $\{t..\} \subseteq \text{existence-ivl0 } x$

$\langle proof \rangle$

lemma $trapped\text{-sol-left}$:

— TODO: when building on `afp-devel`: <https://bitbucket.org/isa-afp/afp-devel/commits/0c3edf9248d5389197f248c723b625c419e4d3eb>

assumes $compact\ K\ K \subseteq X$

assumes $x \in X\ trapped\text{-backward } x K$

shows $\{..0\} \subseteq \text{existence-ivl0 } x$

$\langle proof \rangle$

lemma $trapped\text{-sol-left-gen}$:

assumes $compact\ K\ K \subseteq X$

assumes $t \in \text{existence-ivl0 } x\ trapped\text{-backward } (flow0\ x\ t)\ K$

shows $\{..t\} \subseteq \text{existence-ivl0 } x$

$\langle proof \rangle$

lemma *trapped-sol*:

assumes *compact* $K \subseteq X$
assumes $x \in X$ *trapped* $x \subseteq K$
shows *existence-ivl* $x = UNIV$
(*proof*)

lemma *regular-locally-noteq*:— TODO: should be true in *ll-on-open-it*

assumes $x \in X$ $f x \neq 0$
shows *eventually* $(\lambda t. \text{flow0 } x t \neq x)$ (at 0)
(*proof*)

lemma *compact-max-time-flow-in-closed*:

assumes *closed* M **and** *t-ex*: $t \in \text{existence-ivl} 0 x$
shows *compact* $\{s \in \{0..t\}. \text{flow0 } x ' \{0..s\} \subseteq M\}$ (**is compact** ? C)
(*proof*)

lemma *flow-in-closed-max-timeE*:

assumes *closed* M $t \in \text{existence-ivl} 0 x$ $0 \leq t$ $x \in M$
obtains T **where** $0 \leq T$ $T \leq t$ $\text{flow0 } x ' \{0..T\} \subseteq M$
 $\wedge s'. 0 \leq s' \implies s' \leq t \implies \text{flow0 } x ' \{0..s'\} \subseteq M \implies s' \leq T$
(*proof*)

lemma *flow-leaves-closed-at-frontierE*:

assumes *closed* M **and** *t-ex*: $t \in \text{existence-ivl} 0 x$ **and** $0 \leq t$ $x \in M$ $\text{flow0 } x t \notin M$
obtains s **where** $0 \leq s$ $s < t$ $\text{flow0 } x ' \{0..s\} \subseteq M$
 $\text{flow0 } x s \in \text{frontier } M$
 $\exists_F s'$ *in at-right* $s. \text{flow0 } x s' \notin M$
(*proof*)

2.5 Connectedness

lemma *fcontX*:

shows *continuous-on* X f
(*proof*)

lemma *fcontx*:

assumes $x \in X$
shows *continuous* (at x) f
(*proof*)

lemma *continuous-at-imp-cball*:

assumes *continuous* (at x) g
assumes $g x > (0::\text{real})$
obtains r **where** $r > 0$ $\forall y \in \text{cball } x r. g y > 0$
(*proof*)

flow0 is *path-connected*

lemma *flow0-path-connected-time*:
assumes $ts \subseteq \text{existence-ivl0 } x \text{ path-connected } ts$
shows $\text{path-connected } (\text{flow0 } x \text{ ' } ts)$
 $\langle \text{proof} \rangle$

lemma *flow0-path-connected*:
assumes $\text{path-connected } D$
 $\text{path-connected } ts$
 $\bigwedge x. x \in D \implies ts \subseteq \text{existence-ivl0 } x$
shows $\text{path-connected } ((\lambda(x, y). \text{flow0 } x y) \text{ ' } (D \times ts))$
 $\langle \text{proof} \rangle$

end

2.6 Return Time and Implicit Function Theorem

context *c1-on-open-euclidean* **begin**

lemma *flow-implicit-function*:
— TODO: generalization of $\llbracket \text{returns-to } \{x \in ?S. ?s x = 0\} ?x; \text{closed } ?S; \bigwedge x. (?s \text{ has-derivative } \text{blinfun-apply } (?Ds x)) \text{ (at } x); \text{isCont } ?Ds (\text{poincare-map } \{x \in ?S. ?s x = 0\} ?x); \text{blinfun-apply } (?Ds (\text{poincare-map } \{x \in ?S. ?s x = 0\} ?x)) (f (\text{poincare-map } \{x \in ?S. ?s x = 0\} ?x)) \neq 0; \bigwedge u e. \llbracket ?s (\text{flow0 } ?x (u ?x)) = 0; u ?x = \text{return-time } \{x \in ?S. ?s x = 0\} ?x; \bigwedge y. y \in \text{cball } ?x e \implies ?s (\text{flow0 } y (u y)) = 0; \text{continuous-on } (\text{cball } ?x e) u; (\lambda t. (t, u t)) \text{ ' } \text{cball } ?x e \subseteq \text{Sigma } X \text{ existence-ivl0}; 0 < e; (u \text{ has-derivative } \text{blinfun-apply } (- \text{blinfun-scaleR-left } (\text{inverse } (\text{blinfun-apply } (?Ds (\text{poincare-map } \{x \in ?S. ?s x = 0\} ?x)) (f (\text{poincare-map } \{x \in ?S. ?s x = 0\} ?x)))))) \text{ o}_L (?Ds (\text{poincare-map } \{x \in ?S. ?s x = 0\} ?x) \text{ o}_L \text{flowderiv } ?x (\text{return-time } \{x \in ?S. ?s x = 0\} ?x)) \text{ o}_L \text{embed1-blinfun}) \rrbracket \implies ?thesis \rrbracket \implies ?thesis!$
fixes $s::'a::\text{euclidean-space} \Rightarrow \text{real}$ **and** $S::'a \text{ set}$
assumes $t: t \in \text{existence-ivl0 } x$ **and** $x: x \in X$ **and** $st: s (\text{flow0 } x t) = 0$
assumes $Ds: \bigwedge x. (s \text{ has-derivative } \text{blinfun-apply } (Ds x)) \text{ (at } x)$
assumes $DsC: \text{isCont } Ds (\text{flow0 } x t)$
assumes $nz: Ds (\text{flow0 } x t) (f (\text{flow0 } x t)) \neq 0$
obtains $u e$
where $s (\text{flow0 } x (u x)) = 0$
 $u x = t$
 $(\bigwedge y. y \in \text{cball } x e \implies s (\text{flow0 } y (u y)) = 0)$
 $\text{continuous-on } (\text{cball } x e) u$
 $(\lambda t. (t, u t)) \text{ ' } \text{cball } x e \subseteq \text{Sigma } X \text{ existence-ivl0}$
 $0 < e (u \text{ has-derivative } (- \text{blinfun-scaleR-left } (\text{inverse } (\text{blinfun-apply } (Ds (\text{flow0 } x t)) (f (\text{flow0 } x t)))))) \text{ o}_L$
 $(Ds (\text{flow0 } x t) \text{ o}_L \text{flowderiv } x t) \text{ o}_L \text{embed1-blinfun}) \text{ (at } x)$
 $\langle \text{proof} \rangle$

lemma *flow-implicit-function-at*:
fixes $s::'a::\text{euclidean-space} \Rightarrow \text{real}$ **and** $S::'a \text{ set}$
assumes $x: x \in X$ **and** $st: s x = 0$
assumes $Ds: \bigwedge x. (s \text{ has-derivative } \text{blinfun-apply } (Ds x)) \text{ (at } x)$

assumes $DsC: isCont\ Ds\ x$
assumes $nz: Ds\ x\ (f\ x) \neq 0$
assumes $pos: e > 0$
obtains $u\ d$
where
 $0 < d$
 $u\ x = 0$
 $\bigwedge y. y \in cball\ x\ d \implies s\ (flow0\ y\ (u\ y)) = 0$
 $\bigwedge y. y \in cball\ x\ d \implies |u\ y| < e$
 $\bigwedge y. y \in cball\ x\ d \implies u\ y \in existence-ivl0\ y$
 $continuous-on\ (cball\ x\ d)\ u$
 $(u\ has-derivative\ -Ds\ x\ /_R\ (Ds\ x)\ (f\ x))\ (at\ x)$
 $\langle proof \rangle$

lemma *returns-to-implicit-function-gen:*

— TODO: generalizes proof of $\llbracket returns-to\ \{x \in ?S. ?s\ x = 0\}\ ?x; closed\ ?S; \bigwedge x. (?s\ has-derivative\ blinfun-apply\ (?Ds\ x))\ (at\ x); isCont\ ?Ds\ (poincare-map\ \{x \in ?S. ?s\ x = 0\}\ ?x); blinfun-apply\ (?Ds\ (poincare-map\ \{x \in ?S. ?s\ x = 0\}\ ?x))\ (f\ (poincare-map\ \{x \in ?S. ?s\ x = 0\}\ ?x)) \neq 0; \bigwedge u\ e. \llbracket ?s\ (flow0\ ?x\ (u\ ?x)) = 0; u\ ?x = return-time\ \{x \in ?S. ?s\ x = 0\}\ ?x; \bigwedge y. y \in cball\ ?x\ e \implies ?s\ (flow0\ y\ (u\ y)) = 0; continuous-on\ (cball\ ?x\ e)\ u; (\lambda t. (t, u\ t))\ 'cball\ ?x\ e \subseteq Sigma\ X\ existence-ivl0; 0 < e; (u\ has-derivative\ blinfun-apply\ (-\ blinfun-scaleR-left\ (inverse\ (blinfun-apply\ (?Ds\ (poincare-map\ \{x \in ?S. ?s\ x = 0\}\ ?x))\ (f\ (poincare-map\ \{x \in ?S. ?s\ x = 0\}\ ?x))))\ o_L\ (?Ds\ (poincare-map\ \{x \in ?S. ?s\ x = 0\}\ ?x)\ o_L\ flowderiv\ ?x\ (return-time\ \{x \in ?S. ?s\ x = 0\}\ ?x))\ o_L\ embed1-blinfun)\ (at\ ?x) \rrbracket \implies ?thesis \rrbracket \implies ?thesis!$

fixes $s::'a::euclidean-space \Rightarrow real$
assumes $rt: returns-to\ \{x \in S. s\ x = 0\}\ x\ (is\ returns-to\ ?P\ x)$
assumes $cS: closed\ S$
assumes $Ds: \bigwedge x. (s\ has-derivative\ blinfun-apply\ (Ds\ x))\ (at\ x)$
 $isCont\ Ds\ (poincare-map\ ?P\ x)$
 $Ds\ (poincare-map\ ?P\ x)\ (f\ (poincare-map\ ?P\ x)) \neq 0$
obtains $u\ e$
where $s\ (flow0\ x\ (u\ x)) = 0$
 $u\ x = return-time\ ?P\ x$
 $(\bigwedge y. y \in cball\ x\ e \implies s\ (flow0\ y\ (u\ y)) = 0)$
 $continuous-on\ (cball\ x\ e)\ u$
 $(\lambda t. (t, u\ t))\ 'cball\ x\ e \subseteq Sigma\ X\ existence-ivl0$
 $0 < e\ (u\ has-derivative\ (-\ blinfun-scaleR-left\ (inverse\ (blinfun-apply\ (Ds\ (poincare-map\ ?P\ x))\ (f\ (poincare-map\ ?P\ x))))\ o_L\ (Ds\ (poincare-map\ ?P\ x)\ o_L\ flowderiv\ x\ (return-time\ ?P\ x))\ o_L\ embed1-blinfun)\ (at\ x)$
 $\langle proof \rangle$

c.f. Perko Section 3.7 Lemma 2 part 1.

lemma *flow-transversal-surface-finite-intersections:*

fixes $s::'a \Rightarrow 'b::real-normed-vector$
and $Ds::'a \Rightarrow 'a \Rightarrow_L 'b$
assumes $closed\ S$

assumes $\bigwedge x. (s \text{ has-derivative } (Ds \ x)) \ (at \ x)$
assumes $\bigwedge x. x \in S \implies s \ x = 0 \implies Ds \ x \ (f \ x) \neq 0$
assumes $a \leq b \ \{a \ .. \ b\} \subseteq \text{existence-ivl0 } x$
shows $\text{finite } \{t \in \{a..b\}. \text{flow0 } x \ t \in \{x \in S. s \ x = 0\}\}$
 — TODO: define notion of (compact/closed)-(continuous/differentiable/C1)-
 surface?
 <proof>

lemma *uniform-limit-flow0-state*:— TODO: is that something more general?
assumes *compact* C
assumes $C \subseteq X$
shows *uniform-limit* $C \ (\lambda s \ x. \text{flow0 } x \ s) \ (\lambda x. \text{flow0 } x \ 0) \ (at \ 0)$
 <proof>

end

2.7 Fixpoints

context *auto-ll-on-open* **begin**

lemma *fixpoint-sol*:
assumes $x \in X \ f \ x = 0$
shows *existence-ivl0* $x = \text{UNIV flow0 } x \ t = x$
 <proof>

end

end

3 Invariance

theory *Invariance*
imports *ODE-Misc*
begin

context *auto-ll-on-open* **begin**

definition *invariant* $M \longleftrightarrow (\forall x \in M. \text{trapped } x \ M)$

definition *positively-invariant* $M \longleftrightarrow (\forall x \in M. \text{trapped-forward } x \ M)$

definition *negatively-invariant* $M \longleftrightarrow (\forall x \in M. \text{trapped-backward } x \ M)$

lemma *positively-invariant-iff*:
positively-invariant $M \longleftrightarrow$
 $(\bigcup x \in M. \text{flow0 } x \ '(\text{existence-ivl0 } x \ \cap \ \{0..\})) \subseteq M$
 <proof>

lemma *negatively-invariant-iff*:

negatively-invariant $M \longleftrightarrow$
 $(\bigcup x \in M. \text{flow0 } x \text{ ' (existence-ivl0 } x \cap \{..0\})) \subseteq M$
 ⟨proof⟩

lemma *invariant-iff-pos-and-neg-invariant*:
invariant $M \longleftrightarrow$ *positively-invariant* $M \wedge$ *negatively-invariant* M
 ⟨proof⟩

lemma *invariant-iff*:
invariant $M \longleftrightarrow (\bigcup x \in M. \text{flow0 } x \text{ ' (existence-ivl0 } x)) \subseteq M$
 ⟨proof⟩

lemma *positively-invariant-restrict-dom*: *positively-invariant* $M =$ *positively-invariant*
 $(M \cap X)$
 ⟨proof⟩

lemma *negatively-invariant-restrict-dom*: *negatively-invariant* $M =$ *negatively-invariant*
 $(M \cap X)$
 ⟨proof⟩

lemma *invariant-restrict-dom*: *invariant* $M =$ *invariant* $(M \cap X)$
 ⟨proof⟩

end context *auto-ll-on-open begin*

lemma *positively-invariant-eq-rev*: *positively-invariant* $M =$ *rev.negatively-invariant*
 M
 ⟨proof⟩

lemma *negatively-invariant-eq-rev*: *negatively-invariant* $M =$ *rev.positively-invariant*
 M
 ⟨proof⟩

lemma *invariant-eq-rev*: *invariant* $M =$ *rev.invariant* M
 ⟨proof⟩

lemma *negatively-invariant-complI*: *negatively-invariant* $(X - M)$ **if** *positively-invariant*
 M
 ⟨proof⟩

end context *auto-ll-on-open begin*

lemma *negatively-invariant-complD*: *positively-invariant* M **if** *negatively-invariant*
 $(X - M)$
 ⟨proof⟩

lemma *pos-invariant-iff-compl-neg-invariant*: *positively-invariant* $M \longleftrightarrow$ *negatively-invariant*
 $(X - M)$

<proof>

lemma *neg-invariant-iff-compl-pos-invariant:*

shows *negatively-invariant* $M \longleftrightarrow$ *positively-invariant* $(X - M)$

<proof>

lemma *invariant-iff-compl-invariant:*

shows *invariant* $M \longleftrightarrow$ *invariant* $(X - M)$

<proof>

lemma *invariant-iff-pos-invariant-and-compl-pos-invariant:*

shows *invariant* $M \longleftrightarrow$ *positively-invariant* $M \wedge$ *positively-invariant* $(X - M)$

<proof>

end

3.1 Tools for proving invariance

context *auto-ll-on-open begin*

lemma *positively-invariant-left-inter:*

assumes *positively-invariant* C

assumes $\forall x \in C \cap D.$ *trapped-forward* $x D$

shows *positively-invariant* $(C \cap D)$

<proof>

lemma *trapped-forward-le:*

fixes $V :: 'a \Rightarrow \text{real}$

assumes $V x \leq 0$

assumes *contg: continuous-on* $(\text{flow0 } x \text{ ' (existence-ivl0 } x \cap \{0..\})$ g

assumes $\bigwedge x. (V \text{ has-derivative } V' x) \text{ (at } x)$

assumes $\bigwedge s. s \in \text{existence-ivl0 } x \cap \{0..\} \implies V' (\text{flow0 } x s) (f (\text{flow0 } x s)) \leq g$
 $(\text{flow0 } x s) * V (\text{flow0 } x s)$

shows *trapped-forward* $x \{x. V x \leq 0\}$

<proof>

lemma *positively-invariant-le-domain:*

fixes $V :: 'a \Rightarrow \text{real}$

assumes *positively-invariant* D

assumes *contg: continuous-on* $D g$

assumes $\bigwedge x. (V \text{ has-derivative } V' x) \text{ (at } x)$

assumes $\bigwedge s. s \in D \implies V' s (f s) \leq g s * V s$

shows *positively-invariant* $(D \cap \{x. V x \leq 0\})$

<proof>

lemma *positively-invariant-barrier-domain:*

fixes $V :: 'a \Rightarrow \text{real}$

assumes *positively-invariant* D

assumes $\bigwedge x. (V \text{ has-derivative } V' x) \text{ (at } x)$

assumes *continuous-on* D $(\lambda x. V' x (f x))$
assumes $\bigwedge s. s \in D \implies V s = 0 \implies V' s (f s) < 0$
shows *positively-invariant* $(D \cap \{x. V x \leq 0\})$
 <proof>

lemma *positively-invariant-UNIV*:
shows *positively-invariant* *UNIV*
 <proof>

lemma *positively-invariant-conj*:
assumes *positively-invariant* C
assumes *positively-invariant* D
shows *positively-invariant* $(C \cap D)$
 <proof>

lemma *positively-invariant-le*:
fixes $V :: 'a \Rightarrow \text{real}$
assumes *contg: continuous-on UNIV* g
assumes $\bigwedge x. (V \text{ has-derivative } V' x) (at x)$
assumes $\bigwedge s. V' s (f s) \leq g s * V s$
shows *positively-invariant* $\{x. V x \leq 0\}$
 <proof>

lemma *positively-invariant-barrier*:
fixes $V :: 'a \Rightarrow \text{real}$
assumes $\bigwedge x. (V \text{ has-derivative } V' x) (at x)$
assumes *continuous-on UNIV* $(\lambda x. V' x (f x))$
assumes $\bigwedge s. V s = 0 \implies V' s (f s) < 0$
shows *positively-invariant* $\{x. V x \leq 0\}$
 <proof>

end

end

4 Limit Sets

theory *Limit-Set*
imports *Invariance*
begin

context *auto-ll-on-open* **begin**

Positive limit point, assuming $\{0..\} \subseteq \text{existence-ivl0 } x$

definition *ω -limit-point* $x p \longleftrightarrow$
 $\{0..\} \subseteq \text{existence-ivl0 } x \wedge$
 $(\exists s. s \longrightarrow \infty \wedge (\text{flow0 } x \circ s) \longrightarrow p)$

Also called the ω -limit set of x

definition ω -limit-set $x = \{p. \omega\text{-limit-point } x \ p\}$

definition α -limit-point $x \ p \longleftrightarrow$
 $\{..0\} \subseteq \text{existence-ivl0 } x \wedge$
 $(\exists s. s \longrightarrow -\infty \wedge (\text{flow0 } x \circ s) \longrightarrow p)$

Also called the α -limit set of x

definition α -limit-set $x =$
 $\{p. \alpha\text{-limit-point } x \ p\}$

end context *auto-ll-on-open* **begin**

lemma α -limit-point-eq-rev: $\alpha\text{-limit-point } x \ p = \text{rev.}\omega\text{-limit-point } x \ p$
<proof>

lemma α -limit-set-eq-rev: $\alpha\text{-limit-set } x = \text{rev.}\omega\text{-limit-set } x$
<proof>

lemma ω -limit-pointE:
assumes $\omega\text{-limit-point } x \ p$
obtains s **where**
filterlim s *at-top* *sequentially*
 $(\text{flow0 } x \circ s) \longrightarrow p$
 $\forall n. b \leq s \ n$
<proof>

lemma ω -limit-set-eq:
assumes $\{0..\} \subseteq \text{existence-ivl0 } x$
shows $\omega\text{-limit-set } x = (\text{INF } \tau \in \{0..\}. \text{closure } (\text{flow0 } x \ ' \ \{\tau..\}))$
<proof>

lemma ω -limit-set-empty:
assumes $\neg (\{0..\} \subseteq \text{existence-ivl0 } x)$
shows $\omega\text{-limit-set } x = \{\}$
<proof>

lemma ω -limit-set-closed: *closed* ($\omega\text{-limit-set } x$)
<proof>

lemma ω -limit-set-positively-invariant:
shows *positively-invariant* ($\omega\text{-limit-set } x$)
<proof>

lemma ω -limit-set-invariant:
shows *invariant* ($\omega\text{-limit-set } x$)
<proof>

end context *auto-ll-on-open* **begin**

lemma α -limit-set-eq:
assumes $\{..0\} \subseteq \text{existence-ivl0 } x$
shows $\alpha\text{-limit-set } x = (\text{INF } \tau \in \{..0\}. \text{closure } (\text{flow0 } x \text{ ' } \{..\tau\}))$
 $\langle \text{proof} \rangle$

lemma α -limit-set-closed:
shows $\text{closed } (\alpha\text{-limit-set } x)$
 $\langle \text{proof} \rangle$

lemma α -limit-set-positively-invariant:
shows $\text{negatively-invariant } (\alpha\text{-limit-set } x)$
 $\langle \text{proof} \rangle$

lemma α -limit-set-invariant:
shows $\text{invariant } (\alpha\text{-limit-set } x)$
 $\langle \text{proof} \rangle$

Fundamental properties of the positive limit set

context
fixes $x K$
assumes $K: \text{compact } K \subseteq X$
assumes $x: x \in X \text{ trapped-forward } x K$
begin

Bunch of facts for what's to come

private lemma *props*:
shows $\{0..\} \subseteq \text{existence-ivl0 } x \text{ seq-compact } K$
 $\langle \text{proof} \rangle$ **lemma** *flowimg*:
shows $\text{flow0 } x \text{ ' } (\text{existence-ivl0 } x \cap \{0..\}) = \text{flow0 } x \text{ ' } \{0..\}$
 $\langle \text{proof} \rangle$

lemma ω -limit-set-in-compact-subset:
shows $\omega\text{-limit-set } x \subseteq K$
 $\langle \text{proof} \rangle$

lemma ω -limit-set-in-compact-compact:
shows $\text{compact } (\omega\text{-limit-set } x)$
 $\langle \text{proof} \rangle$

lemma ω -limit-set-in-compact-nonempty:
shows $\omega\text{-limit-set } x \neq \{\}$
 $\langle \text{proof} \rangle$

lemma ω -limit-set-in-compact-existence:
shows $\bigwedge y. y \in \omega\text{-limit-set } x \implies \text{existence-ivl0 } y = \text{UNIV}$
 $\langle \text{proof} \rangle$

lemma ω -limit-set-in-compact-tendsto:

shows $((\lambda t. \text{infdist} (\text{flow0 } x \ t) (\omega\text{-limit-set } x)) \longrightarrow 0)$ *at-top*
<proof>

lemma *ω -limit-set-in-compact-connected:*

shows *connected* $(\omega\text{-limit-set } x)$
<proof>

lemma *ω -limit-set-in-compact- ω -limit-set-contained:*

shows $\forall y \in \omega\text{-limit-set } x. \omega\text{-limit-set } y \subseteq \omega\text{-limit-set } x$
<proof>

lemma *ω -limit-set-in-compact- α -limit-set-contained:*

assumes *zpx: $z \in \omega\text{-limit-set } x$*
shows *$\alpha\text{-limit-set } z \subseteq \omega\text{-limit-set } x$*
<proof>

end

Fundamental properties of the negative limit set

end context *auto-ll-on-open* **begin**

context

fixes *$x \ K$*

assumes *$x: x \in X$ *trapped-backward* $x \ K$*

assumes *$K: \text{compact } K \ K \subseteq X$*

begin

private lemma *$xrev: x \in X$ *rev.trapped-forward* $x \ K$*

<proof>

lemma *α -limit-set-in-compact-subset: $\alpha\text{-limit-set } x \subseteq K$*

and *α -limit-set-in-compact-compact: *compact* $(\alpha\text{-limit-set } x)$*

and *α -limit-set-in-compact-nonempty: $\alpha\text{-limit-set } x \neq \{\}$*

and *α -limit-set-in-compact-connected: *connected* $(\alpha\text{-limit-set } x)$*

and *α -limit-set-in-compact- α -limit-set-contained:*

$\forall y \in \alpha\text{-limit-set } x. \alpha\text{-limit-set } y \subseteq \alpha\text{-limit-set } x$

and *α -limit-set-in-compact-tendsto: $((\lambda t. \text{infdist} (\text{flow0 } x \ t) (\alpha\text{-limit-set } x)) \longrightarrow 0)$ *at-bot**

<proof>

lemma *α -limit-set-in-compact-existence:*

shows $\bigwedge y. y \in \alpha\text{-limit-set } x \implies \text{existence-ivl0 } y = \text{UNIV}$

<proof>

end

end

end

5 Periodic Orbits

theory *Periodic-Orbit*

imports

Ordinary-Differential-Equations.ODE-Analysis

Analysis-Misc

ODE-Misc

Limit-Set

begin

Definition of closed and periodic orbits and their associated properties

context *auto-ll-on-open*

begin

TODO: not sure if the "closed orbit" terminology is standard Closed orbits have some non-zero recurrence time T where the flow returns to the initial state The period of a closed orbit is the infimum of all positive recurrence times Periodic orbits are the subset of closed orbits where the period is non-zero

definition *closed-orbit* $x \longleftrightarrow$

$(\exists T \in \text{existence-ivl0 } x. T \neq 0 \wedge \text{flow0 } x T = x)$

definition *period* $x =$

$\text{Inf } \{T \in \text{existence-ivl0 } x. T > 0 \wedge \text{flow0 } x T = x\}$

definition *periodic-orbit* $x \longleftrightarrow$

$\text{closed-orbit } x \wedge \text{period } x > 0$

lemma *recurrence-time-flip-sign:*

assumes $T \in \text{existence-ivl0 } x \text{ flow0 } x T = x$

shows $-T \in \text{existence-ivl0 } x \text{ flow0 } x (-T) = x$

<proof>

lemma *closed-orbit-recurrence-times-nonempty:*

assumes *closed-orbit* x

shows $\{T \in \text{existence-ivl0 } x. T > 0 \wedge \text{flow0 } x T = x\} \neq \{\}$

<proof>

lemma *closed-orbit-recurrence-times-bdd-below:*

shows *bdd-below* $\{T \in \text{existence-ivl0 } x. T > 0 \wedge \text{flow0 } x T = x\}$

<proof>

lemma *closed-orbit-period-nonneg:*

assumes *closed-orbit* x

shows *period* $x \geq 0$

<proof>

lemma *closed-orbit-in-domain:*

assumes *closed-orbit* x

shows $x \in X$
<proof>

lemma *closed-orbit-global-existence:*

assumes *closed-orbit* x
shows *existence-ivl0* $x = UNIV$
<proof>

lemma *recurrence-time-multiples:*

fixes $n::nat$
assumes $T \in \text{existence-ivl0 } x$ $T \neq 0$ *flow0* x $T = x$
shows $\bigwedge t. \text{flow0 } x (t+T*n) = \text{flow0 } x t$
<proof>

lemma *nasty-arithmetic1:*

fixes $t T::real$
assumes $T > 0$ $t \geq 0$
obtains $q r$ **where** $t = (q::nat) * T + r$ $0 \leq r < T$
<proof>

lemma *nasty-arithmetic2:*

fixes $t T::real$
assumes $T > 0$ $t \leq 0$
obtains $q r$ **where** $t = (q::nat) * (-T) + r$ $0 \leq r < T$
<proof>

lemma *recurrence-time-restricts-compact-flow:*

assumes $T \in \text{existence-ivl0 } x$ $T > 0$ *flow0* x $T = x$
shows *flow0* $x \text{ ' } UNIV = \text{flow0 } x \text{ ' } \{0..T\}$
<proof>

lemma *closed-orbitI:*

assumes $t \neq t'$ $t \in \text{existence-ivl0 } y$ $t' \in \text{existence-ivl0 } y$
assumes *flow0* y $t = \text{flow0 } y t'$
shows *closed-orbit* y
<proof>

lemma *flow0-image-UNIV:*

assumes *existence-ivl0* $x = UNIV$
shows *flow0* $(\text{flow0 } x t) \text{ ' } S = \text{flow0 } x \text{ ' } (\lambda s. s + t) \text{ ' } S$
<proof>

lemma *recurrence-time-restricts-compact-flow':*

assumes $t < t'$ $t \in \text{existence-ivl0 } y$ $t' \in \text{existence-ivl0 } y$
assumes *flow0* y $t = \text{flow0 } y t'$
shows *flow0* $y \text{ ' } UNIV = \text{flow0 } y \text{ ' } \{t..t'\}$
<proof>

lemma *closed-orbitE'*:
assumes *closed-orbit x*
obtains T **where** $T > 0 \wedge t (n::nat). \text{flow0 } x (t+T*n) = \text{flow0 } x t$
 $\langle \text{proof} \rangle$

lemma *closed-orbitE*:
assumes *closed-orbit x*
obtains T **where** $T > 0 \wedge t. \text{flow0 } x (t+T) = \text{flow0 } x t$
 $\langle \text{proof} \rangle$

lemma *closed-orbit-flow-compact*:
assumes *closed-orbit x*
shows $\text{compact}(\text{flow0 } x \text{ ' UNIV})$
 $\langle \text{proof} \rangle$

lemma *fixed-point-imp-closed-orbit-period-zero*:
assumes $x \in X$
assumes $f x = 0$
shows *closed-orbit x period x = 0*
 $\langle \text{proof} \rangle$

lemma *closed-orbit-period-zero-fixed-point*:
assumes *closed-orbit x period x = 0*
shows $f x = 0$
 $\langle \text{proof} \rangle$

lemma *closed-orbit-subset- ω -limit-set*:
assumes *closed-orbit x*
shows $\text{flow0 } x \text{ ' UNIV} \subseteq \omega\text{-limit-set } x$
 $\langle \text{proof} \rangle$

lemma *closed-orbit- ω -limit-set*:
assumes *closed-orbit x*
shows $\text{flow0 } x \text{ ' UNIV} = \omega\text{-limit-set } x$
 $\langle \text{proof} \rangle$

lemma *flow0-inj-on*:
assumes $t \leq t'$
assumes $\{t..t'\} \subseteq \text{existence-ivl0 } x$
assumes $\wedge s. t < s \implies s \leq t' \implies \text{flow0 } x s \neq \text{flow0 } x t$
shows *inj-on (flow0 x) {t..t'}*
 $\langle \text{proof} \rangle$

lemma *finite- ω -limit-set-in-compact-imp-unique-fixed-point*:
assumes $\text{compact } K \ K \subseteq X$
assumes $x \in X \ \text{trapped-forward } x \ K$
assumes *finite (ω -limit-set x)*
obtains y **where** $\omega\text{-limit-set } x = \{y\} \ f y = 0$

<proof>

lemma *closed-orbit-periodic:*

assumes *closed-orbit* $x \neq 0$

shows *periodic-orbit* x

<proof>

lemma *periodic-orbitI:*

assumes $t \neq t'$ $t \in \text{existence-ivl0 } y$ $t' \in \text{existence-ivl0 } y$

assumes $\text{flow0 } y \ t = \text{flow0 } y \ t'$

assumes $f \ y \neq 0$

shows *periodic-orbit* y

<proof>

lemma *periodic-orbit-recurrence-times-closed:*

assumes *periodic-orbit* x

shows $\text{closed } \{T \in \text{existence-ivl0 } x. T > 0 \wedge \text{flow0 } x \ T = x\}$

<proof>

lemma *periodic-orbit-period:*

assumes *periodic-orbit* x

shows $\text{period } x > 0 \ \text{flow0 } x \ (\text{period } x) = x$

<proof>

lemma *closed-orbit-flow0:*

assumes *closed-orbit* x

shows *closed-orbit* $(\text{flow0 } x \ t)$

<proof>

lemma *periodic-orbit-imp-flow0-regular:*

assumes *periodic-orbit* x

shows $f \ (\text{flow0 } x \ t) \neq 0$

<proof>

lemma *fixed-point-imp- ω -limit-set:*

assumes $x \in X$ $f \ x = 0$

shows $\omega\text{-limit-set } x = \{x\}$

<proof>

end

context *auto-ll-on-open* **begin**

lemma *closed-orbit-eq-rev:* $\text{closed-orbit } x = \text{rev.closed-orbit } x$

<proof>

lemma *closed-orbit- α -limit-set:*

assumes *closed-orbit* x

shows $\text{flow0 } x \ ' \ \text{UNIV} = \alpha\text{-limit-set } x$

<proof>

lemma *fixed-point-imp- α -limit-set:*

assumes $x \in X$ $f x = 0$

shows α -limit-set $x = \{x\}$

<proof>

lemma *finite- α -limit-set-in-compact-imp-unique-fixed-point:*

assumes compact K $K \subseteq X$

assumes $x \in X$ *trapped-backward* $x K$

assumes *finite* (α -limit-set x)

obtains y **where** α -limit-set $x = \{y\}$ $f y = 0$

<proof>

end

end

6 Poincare Bendixson Theory

theory *Poincare-Bendixson*

imports

Ordinary-Differential-Equations.ODE-Analysis

Analysis-Misc ODE-Misc Periodic-Orbit

begin

6.1 Flow to Path

context *auto-ll-on-open* **begin**

definition *flow-to-path* $x t t' = \text{flow0 } x \circ \text{linepath } t t'$

lemma *pathstart-flow-to-path[simp]:*

shows *pathstart* (*flow-to-path* $x t t'$) = *flow0* $x t$

<proof>

lemma *pathfinish-flow-to-path[simp]:*

shows *pathfinish* (*flow-to-path* $x t t'$) = *flow0* $x t'$

<proof>

lemma *flow-to-path-unfold:*

shows *flow-to-path* $x t t' s = \text{flow0 } x ((1 - s) * t + s * t')$

<proof>

lemma *subpath0-flow-to-path:*

shows (*subpath* 0 u (*flow-to-path* $x t t'$)) = *flow-to-path* $x t (t + u*(t'-t))$

<proof>

lemma *path-image-flow-to-path[simp]:*

assumes $t \leq t'$
shows $\text{path-image } (\text{flow-to-path } x \ t \ t') = \text{flow0 } x \ \{t..t'\}$
 $\langle \text{proof} \rangle$

lemma $\text{flow-to-path-image0-right-open}[simp]$:
assumes $t < t'$
shows $\text{flow-to-path } x \ t \ t' \ \{0..<1\} = \text{flow0 } x \ \{t..<t'\}$
 $\langle \text{proof} \rangle$

lemma flow-to-path-path :
assumes $t \leq t'$
assumes $\{t..t'\} \subseteq \text{existence-ivl0 } x$
shows $\text{path } (\text{flow-to-path } x \ t \ t')$
 $\langle \text{proof} \rangle$

lemma flow-to-path-arc :
assumes $t \leq t'$
assumes $\{t..t'\} \subseteq \text{existence-ivl0 } x$
assumes $\forall s \in \{t<..<t'\}. \text{flow0 } x \ s \neq \text{flow0 } x \ t$
assumes $\text{flow0 } x \ t \neq \text{flow0 } x \ t'$
shows $\text{arc } (\text{flow-to-path } x \ t \ t')$
 $\langle \text{proof} \rangle$

end

locale $c1\text{-on-open-}R2 = c1\text{-on-open-euclidean } f \ f' \ X \ \text{for } f::'a::\text{executable-euclidean-space}$
 \Rightarrow - **and** f' **and** X +
assumes $\text{dim2}: \text{DIM}('a) = 2$
begin

6.2 2D Line segments

Line segments are specified by two endpoints The closed line segment from x to y is given by the set $x..y$ and $x<..<y$ for the open segment

Rotates a vector clockwise 90 degrees

definition $\text{rot } (v::'a) = (\text{eucl-of-list } [nth\text{-eucl } v \ 1, -nth\text{-eucl } v \ 0]::'a)$

lemma $\text{exhaust2-nat}: (\forall i < (2::nat). P \ i) \longleftrightarrow P \ 0 \ \wedge \ P \ 1$
 $\langle \text{proof} \rangle$

lemma $\text{sum2-nat}: (\sum i < (2::nat). P \ i) = P \ 0 \ + \ P \ 1$
 $\langle \text{proof} \rangle$

lemmas $\text{vec-simps} =$
 $\text{eucl-eq-iff}[\text{where } 'a='a] \ \text{dim2 } \text{eucl-of-list-eucl-nth } \text{exhaust2-nat}$
 plus-nth-eucl
 minus-nth-eucl
 uminus-nth-eucl
 scaleR-nth-eucl

inner-nth-eucl
sum2-nat
algebra-simps

lemma *minus-expand*:

shows $(x::'a)-y = (\text{eucl-of-list } [x\$_e0 - y\$_e0, x\$_e1 - y\$_e1])$
 $\langle \text{proof} \rangle$

lemma *dot-ortho[simp]*: $x \cdot \text{rot } x = 0$

$\langle \text{proof} \rangle$

lemma *nrm-dot*:

shows $((x::'a)-y) \cdot (\text{rot } (x-y)) = 0$
 $\langle \text{proof} \rangle$

lemma *nrm-reverse*: $a \cdot (\text{rot } (x-y)) = -a \cdot (\text{rot } (y-x))$ **for** $x\ y::'a$

$\langle \text{proof} \rangle$

lemma *norm-rot*: $\text{norm } (\text{rot } v) = \text{norm } v$ **for** $v::'a$

$\langle \text{proof} \rangle$

lemma *rot-rot[simp]*:

shows $\text{rot } (\text{rot } v) = -v$
 $\langle \text{proof} \rangle$

lemma *rot-scaleR[simp]*:

shows $\text{rot } (u *_R v) = u *_R (\text{rot } v)$
 $\langle \text{proof} \rangle$

lemma *rot-0[simp]*: $\text{rot } 0 = 0$

$\langle \text{proof} \rangle$

lemma *rot-eq-0-iff[simp]*: $\text{rot } x = 0 \longleftrightarrow x = 0$

$\langle \text{proof} \rangle$

lemma *in-segment-inner-rot*:

$(x - a) \cdot \text{rot } (b - a) = 0$
if $x \in \{a \dashv\vdash b\}$

$\langle \text{proof} \rangle$

lemma *inner-rot-in-segment*:

$x \in \text{range } (\lambda u. a + u *_R (b - a))$
if $(x - a) \cdot \text{rot } (b - a) = 0$ $a \neq b$

$\langle \text{proof} \rangle$

lemma *in-open-segment-iff-rot*:

$x \in \{a < \dashv\vdash < b\} \longleftrightarrow (x - a) \cdot \text{rot } (b - a) = 0 \wedge x \cdot (b - a) \in \{a \cdot (b - a) < .. < b \cdot (b - a)\}$
if $a \neq b$

<proof>

lemma *in-open-segment-rotD:*

$x \in \{a <--<b\} \implies (x - a) \cdot \text{rot}(b - a) = 0 \wedge x \cdot (b - a) \in \{a \cdot (b - a) <..<b \cdot (b - a)\}$
<proof>

lemma *in-closed-segment-iff-rot:*

$x \in \{a--b\} \iff (x - a) \cdot \text{rot}(b - a) = 0 \wedge x \cdot (b - a) \in \{a \cdot (b - a) .. b \cdot (b - a)\}$
if $a \neq b$
<proof>

lemma *in-segment-inner-rot2:*

$(x - y) \cdot \text{rot}(a - b) = 0$
if $x \in \{a--b\} \ y \in \{a--b\}$
<proof>

lemma *closed-segment-surface:*

$a \neq b \implies \{a--b\} = \{x \in \{x. x \cdot (b - a) \in \{a \cdot (b - a) .. b \cdot (b - a)\}\}. (x - a) \cdot \text{rot}(b - a) = 0\}$
<proof>

lemma *rot-diff-commute:* $\text{rot}(b - a) = -\text{rot}(a - b)$

<proof>

6.3 Bijection Real-Complex for Jordan Curve Theorem

definition *complex-of* $(x::'a) = x\$_e 0 + i * x\$_e 1$

definition *real-of* $(x::\text{complex}) = (\text{eucl-of-list } [\text{Re } x, \text{Im } x]::'a)$

lemma *complex-of-linear:*

shows *linear complex-of*
<proof>

lemma *complex-of-bounded-linear:*

shows *bounded-linear complex-of*
<proof>

lemma *real-of-linear:*

shows *linear real-of*
<proof>

lemma *real-of-bounded-linear:*

shows *bounded-linear real-of*
<proof>

lemma *complex-of-real-of:*

$(\text{complex-of} \circ \text{real-of}) = \text{id}$
{proof}

lemma *real-of-complex-of*:
 $(\text{real-of} \circ \text{complex-of}) = \text{id}$
{proof}

lemma *complex-of-bij*:
shows *bij* (*complex-of*)
{proof}

lemma *real-of-bij*:
shows *bij* (*real-of*)
{proof}

lemma *real-of-inj*:
shows *inj* (*real-of*)
{proof}

lemma *Jordan-curve-R2*:
fixes $c :: \text{real} \Rightarrow 'a$
assumes *simple-path* c *pathfinish* $c = \text{pathstart } c$
obtains *inside* *outside* **where**
 $\text{inside} \neq \{\}$ *open* *inside* *connected* *inside*
 $\text{outside} \neq \{\}$ *open* *outside* *connected* *outside*
 $\text{bounded } \text{inside} \neg \text{bounded } \text{outside}$
 $\text{inside} \cap \text{outside} = \{\}$
 $\text{inside} \cup \text{outside} = - \text{path-image } c$
 $\text{frontier } \text{inside} = \text{path-image } c$
 $\text{frontier } \text{outside} = \text{path-image } c$
{proof}

corollary *Jordan-inside-outside-R2*:
fixes $c :: \text{real} \Rightarrow 'a$
assumes *simple-path* c *pathfinish* $c = \text{pathstart } c$
shows $\text{inside}(\text{path-image } c) \neq \{\}$ \wedge
 $\text{open}(\text{inside}(\text{path-image } c)) \wedge$
 $\text{connected}(\text{inside}(\text{path-image } c)) \wedge$
 $\text{outside}(\text{path-image } c) \neq \{\}$ \wedge
 $\text{open}(\text{outside}(\text{path-image } c)) \wedge$
 $\text{connected}(\text{outside}(\text{path-image } c)) \wedge$
 $\text{bounded}(\text{inside}(\text{path-image } c)) \wedge$
 $\neg \text{bounded}(\text{outside}(\text{path-image } c)) \wedge$
 $\text{inside}(\text{path-image } c) \cap \text{outside}(\text{path-image } c) = \{\}$ \wedge
 $\text{inside}(\text{path-image } c) \cup \text{outside}(\text{path-image } c) =$
 $- \text{path-image } c \wedge$
 $\text{frontier}(\text{inside}(\text{path-image } c)) = \text{path-image } c \wedge$
 $\text{frontier}(\text{outside}(\text{path-image } c)) = \text{path-image } c$

$\langle proof \rangle$

lemma *jordan-points-inside-outside*:

fixes $p :: real \Rightarrow 'a$

assumes $0 < e$

assumes *jordan*: $simple\text{-}path\ p\ pathfinish\ p = pathstart\ p$

assumes $x: x \in path\text{-}image\ p$

obtains $y\ z$ **where** $y \in inside\ (path\text{-}image\ p)\ y \in ball\ x\ e$
 $z \in outside\ (path\text{-}image\ p)\ z \in ball\ x\ e$

$\langle proof \rangle$

lemma *eventually-at-open-segment*:

assumes $x \in \{a <--< b\}$

shows $\forall_F y\ in\ at\ x. (y-a) \cdot rot(a-b) = 0 \longrightarrow y \in \{a <--< b\}$

$\langle proof \rangle$

lemma *linepath-ball*:

assumes $x \in \{a <--< b\}$

obtains e **where** $e > 0\ ball\ x\ e \cap \{y. (y-a) \cdot rot(a-b) = 0\} \subseteq \{a <--< b\}$

$\langle proof \rangle$

lemma *linepath-ball-inside-outside*:

fixes $p :: real \Rightarrow 'a$

assumes *jordan*: $simple\text{-}path\ (p\ +++\ linepath\ a\ b)\ pathfinish\ p = a\ pathstart\ p = b$

assumes $x: x \in \{a <--< b\}$

obtains e **where** $e > 0\ ball\ x\ e \cap path\text{-}image\ p = \{\}$

$ball\ x\ e \cap \{y. (y-a) \cdot rot(a-b) > 0\} \subseteq inside\ (path\text{-}image\ (p\ +++\ linepath\ a\ b)) \wedge$

$ball\ x\ e \cap \{y. (y-a) \cdot rot(a-b) < 0\} \subseteq outside\ (path\text{-}image\ (p\ +++\ linepath\ a\ b))$

\vee

$ball\ x\ e \cap \{y. (y-a) \cdot rot(a-b) < 0\} \subseteq inside\ (path\text{-}image\ (p\ +++\ linepath\ a\ b)) \wedge$

$ball\ x\ e \cap \{y. (y-a) \cdot rot(a-b) > 0\} \subseteq outside\ (path\text{-}image\ (p\ +++\ linepath\ a\ b))$

$\langle proof \rangle$

6.4 Transversal Segments

definition *transversal-segment* $a\ b \longleftrightarrow$

$a \neq b \wedge \{a--b\} \subseteq X \wedge$

$(\forall z \in \{a--b\}. f\ z \cdot rot(a-b) \neq 0)$

lemma *transversal-segment-reverse*:

assumes *transversal-segment* $x\ y$

shows *transversal-segment* $y\ x$

$\langle proof \rangle$

lemma transversal-segment-commute: $\text{transversal-segment } x \ y \longleftrightarrow \text{transversal-segment } y \ x$

<proof>

lemma transversal-segment-neg:

assumes $\text{transversal-segment } x \ y$

assumes $w: w \in \{x \dashrightarrow y\}$ **and** $f \ w \cdot \text{rot } (x-y) < 0$

shows $\forall z \in \{x \dashrightarrow y\}. f(z) \cdot \text{rot } (x-y) < 0$

<proof>

lemmas $\text{transversal-segment-sign-less} = \text{transversal-segment-neg}[\text{OF - ends-in-segment}(1)]$

lemma transversal-segment-pos:

assumes $\text{transversal-segment } x \ y$

assumes $w: w \in \{x \dashrightarrow y\}$ $f \ w \cdot \text{rot } (x-y) > 0$

shows $\forall z \in \{x \dashrightarrow y\}. f(z) \cdot \text{rot } (x-y) > 0$

<proof>

lemma transversal-segment-posD:

assumes $\text{transversal-segment } x \ y$

and $\text{pos}: z \in \{x \dashrightarrow y\}$ $f \ z \cdot \text{rot } (x - y) > 0$

shows $x \neq y \ \{x \dashrightarrow y\} \subseteq X \ \wedge z. z \in \{x \dashrightarrow y\} \implies f \ z \cdot \text{rot } (x-y) > 0$

<proof>

lemma transversal-segment-negD:

assumes $\text{transversal-segment } x \ y$

and $\text{pos}: z \in \{x \dashrightarrow y\}$ $f \ z \cdot \text{rot } (x - y) < 0$

shows $x \neq y \ \{x \dashrightarrow y\} \subseteq X \ \wedge z. z \in \{x \dashrightarrow y\} \implies f \ z \cdot \text{rot } (x-y) < 0$

<proof>

lemma transversal-segmentE:

assumes $\text{transversal-segment } x \ y$

obtains $x \neq y \ \{x \dashrightarrow y\} \subseteq X \ \wedge z. z \in \{x \dashrightarrow y\} \implies f \ z \cdot \text{rot } (x - y) > 0$

| $x \neq y \ \{x \dashrightarrow y\} \subseteq X \ \wedge z. z \in \{x \dashrightarrow y\} \implies f \ z \cdot \text{rot } (y - x) > 0$

<proof>

lemma dist-add-vec:

shows $\text{dist } (x + s *_{\mathbb{R}} v) \ x = \text{abs } s * \text{norm } v$

<proof>

lemma transversal-segment-exists:

assumes $x \in X \ f \ x \neq 0$

obtains $a \ b$ **where** $x \in \{a \dashrightarrow b\}$

$\text{transversal-segment } a \ b$

<proof>

Perko Section 3.7 Lemma 2 part 1.

lemma flow-transversal-segment-finite-intersections:

assumes $\text{transversal-segment } a \ b$

assumes $t \leq t' \{t \dots t'\} \subseteq \text{existence-ivl0 } x$
shows $\text{finite } \{s \in \{t..t'\}. \text{flow0 } x \ s \in \{a \dots b\}\}$
 $\langle \text{proof} \rangle$

lemma *transversal-bound-posE*:

assumes *transversal: transversal-segment* $a \ b$
assumes *direction*: $z \in \{a \dots b\} \ f \ z \cdot (\text{rot } (a - b)) > 0$
obtains $d \ B$ **where** $d > 0 \ 0 < B$
 $\bigwedge x \ y. x \in \{a \dots b\} \implies \text{dist } x \ y \leq d \implies f \ y \cdot (\text{rot } (a - b)) \geq B$
 $\langle \text{proof} \rangle$

lemma *transversal-bound-negE*:

assumes *transversal: transversal-segment* $a \ b$
assumes *direction*: $z \in \{a \dots b\} \ f \ z \cdot (\text{rot } (a - b)) < 0$
obtains $d \ B$ **where** $d > 0 \ 0 < B$
 $\bigwedge x \ y. x \in \{a \dots b\} \implies \text{dist } x \ y \leq d \implies f \ y \cdot (\text{rot } (b - a)) \geq B$
 $\langle \text{proof} \rangle$

lemma *leaves-transversal-segmentE*:

assumes *transversal: transversal-segment* $a \ b$
obtains $T \ n$ **where** $T > 0 \ n = a - b \vee n = b - a$
 $\bigwedge x. x \in \{a \dots b\} \implies \{-T..T\} \subseteq \text{existence-ivl0 } x$
 $\bigwedge x \ s. x \in \{a \dots b\} \implies 0 < s \implies s \leq T \implies$
 $(\text{flow0 } x \ s - x) \cdot \text{rot } n > 0$
 $\bigwedge x \ s. x \in \{a \dots b\} \implies -T \leq s \implies s < 0 \implies$
 $(\text{flow0 } x \ s - x) \cdot \text{rot } n < 0$
 $\langle \text{proof} \rangle$

lemma *inner-rot-pos-move-base*: $(x - a) \cdot \text{rot } (a - b) > 0$

if $(x - y) \cdot \text{rot } (a - b) > 0 \ y \in \{a \dots b\}$

$\langle \text{proof} \rangle$

lemma *inner-rot-neg-move-base*: $(x - a) \cdot \text{rot } (a - b) < 0$

if $(x - y) \cdot \text{rot } (a - b) < 0 \ y \in \{a \dots b\}$

$\langle \text{proof} \rangle$

lemma *inner-pos-move-base*: $(x - a) \cdot n > 0$

if $(a - b) \cdot n = 0 \ (x - y) \cdot n > 0 \ y \in \{a \dots b\}$

$\langle \text{proof} \rangle$

lemma *inner-neg-move-base*: $(x - a) \cdot n < 0$

if $(a - b) \cdot n = 0 \ (x - y) \cdot n < 0 \ y \in \{a \dots b\}$

$\langle \text{proof} \rangle$

lemma *rot-same-dir*:

assumes $x1 \in \{a < \dots < b\}$

assumes $x2 \in \{x1 < \dots < b\}$

shows $(y \cdot \text{rot } (a - b) > 0) = (y \cdot \text{rot}(x1 - x2) > 0) \ (y \cdot \text{rot } (a - b) < 0) = (y \cdot$

$\text{rot}(x1-x2) < 0)$
 ⟨proof⟩

6.5 Monotone Step Lemma

lemma *flow0-transversal-segment-monotone-step:*

assumes *transversal-segment* $a\ b$
assumes $t1 \leq t2$ $\{t1..t2\} \subseteq \text{existence-ivl0 } x$
assumes $x1: \text{flow0 } x\ t1 \in \{a<--<b\}$
assumes $x2: \text{flow0 } x\ t2 \in \{\text{flow0 } x\ t1<--<b\}$
assumes $\bigwedge t. t \in \{t1<..
assumes $t > t2$ $t \in \text{existence-ivl0 } x$
shows $\text{flow0 } x\ t \notin \{a<--<\text{flow0 } x\ t2\}$
 ⟨proof⟩$

lemma *open-segment-trichotomy:*

fixes $x\ y\ a\ b::'a$
assumes $x:x \in \{a<--<b\}$
assumes $y:y \in \{a<--<b\}$
shows $x = y \vee y \in \{x<--<b\} \vee y \in \{a<--<x\}$
 ⟨proof⟩

sublocale *rev: c1-on-open-R2* $-f\ -f'$ **rewrites** $-(-f) = f$ **and** $-(-f') = f'$
 ⟨proof⟩

lemma *rev-transversal-segment: rev.transversal-segment* $a\ b = \text{transversal-segment}$
 $a\ b$

⟨proof⟩

lemma *flow0-transversal-segment-monotone-step-reverse:*

assumes *transversal-segment* $a\ b$
assumes $t1 \leq t2$
assumes $\{t1..t2\} \subseteq \text{existence-ivl0 } x$
assumes $x1: \text{flow0 } x\ t1 \in \{a<--<b\}$
assumes $x2: \text{flow0 } x\ t2 \in \{a<--<\text{flow0 } x\ t1\}$
assumes $\bigwedge t. t \in \{t1<..
assumes $t < t1$ $t \in \text{existence-ivl0 } x$
shows $\text{flow0 } x\ t \notin \{a<--<\text{flow0 } x\ t1\}$
 ⟨proof⟩$

lemma *flow0-transversal-segment-monotone-step-reverse2:*

assumes *transversal: transversal-segment* $a\ b$
assumes *time:* $t1 \leq t2$
assumes *exist:* $\{t1..t2\} \subseteq \text{existence-ivl0 } x$
assumes $t1: \text{flow0 } x\ t1 \in \{a<--<b\}$
assumes $t2: \text{flow0 } x\ t2 \in \{\text{flow0 } x\ t1<--<b\}$
assumes $t1t2: \bigwedge t. t \in \{t1<..
assumes $t: t < t1$ $t \in \text{existence-ivl0 } x$
shows $\text{flow0 } x\ t \notin \{\text{flow0 } x\ t1<--<b\}$$

<proof>

lemma *flow0-transversal-segment-monotone-step2:*

assumes *transversal: transversal-segment a b*

assumes *time: t1 ≤ t2*

assumes *exist: {t1..t2} ⊆ existence-ivl0 x*

assumes *t1: flow0 x t1 ∈ {a<--<b}*

assumes *t2: flow0 x t2 ∈ {a<--<flow0 x t1}*

assumes *t1t2: ∧t. t ∈ {t1<..*

shows *∧t. t > t2 ⇒ t ∈ existence-ivl0 x ⇒ flow0 x t ∉ {flow0 x t2<--<b}*

<proof>

lemma *flow0-transversal-segment-monotone:*

assumes *transversal-segment a b*

assumes *t1 ≤ t2*

assumes *{t1..t2} ⊆ existence-ivl0 x*

assumes *x1: flow0 x t1 ∈ {a<--<b}*

assumes *x2: flow0 x t2 ∈ {flow0 x t1<--<b}*

assumes *t > t2 t ∈ existence-ivl0 x*

shows *flow0 x t ∉ {a<--<flow0 x t2}*

<proof>

6.6 Straightening

This lemma uses the implicit function theorem

lemma *cross-time-continuous:*

assumes *transversal-segment a b*

assumes *x ∈ {a<--<b}*

assumes *e > 0*

obtains *d t where d > 0 continuous-on (ball x d) t*

∧y. y ∈ ball x d ⇒ flow0 y (t y) ∈ {a<--<b}

∧y. y ∈ ball x d ⇒ |t y| < e

continuous-on (ball x d) t

t x = 0

<proof>

lemma *ω-limit-crossings:*

assumes *transversal-segment a b*

assumes *pos-ex: {0..} ⊆ existence-ivl0 x*

assumes *ω-limit-point x p*

assumes *p ∈ {a<--<b}*

obtains *s where*

s → ∞

(flow0 x ∘ s) → p

∀F n in sequentially. flow0 x (s n) ∈ {a<--<b} ∧ s n ∈ existence-ivl0 x

<proof>

lemma *filterlim-at-top-tendstoE:*

assumes $e > 0$
assumes *filterlim s at-top sequentially*
assumes $(\text{flow0 } x \circ s) \longrightarrow u$
assumes $\forall_F n$ *in sequentially. P (s n)*
obtains m **where** $m > b$ $P m$ *dist (flow0 x m) u < e*
 <proof>

lemma *open-segment-separate-left:*
fixes $u v x a b::'a$
assumes $u:u \in \{a <--< b\}$
assumes $v:v \in \{u <--< b\}$
assumes $x: \text{dist } x u < \text{dist } u v$ $x \in \{a <--< b\}$
shows $x \in \{a <--< v\}$
 <proof>

lemma *open-segment-separate-right:*
fixes $u v x a b::'a$
assumes $u:u \in \{a <--< b\}$
assumes $v:v \in \{a <--< u\}$
assumes $x: \text{dist } x u < \text{dist } u v$ $x \in \{a <--< b\}$
shows $x \in \{v <--< b\}$
 <proof>

lemma *no-two- ω -limit-points:*
assumes *transversal: transversal-segment a b*
assumes $\{0..\} \subseteq \text{existence-ivl0 } x$
assumes $u: \omega\text{-limit-point } x$ $u \in \{a <--< b\}$
assumes $v: \omega\text{-limit-point } x$ $v \in \{a <--< b\}$
assumes $uv: v \in \{u <--< b\}$
shows *False*
 <proof>

6.7 Unique Intersection

Perko Section 3.7 Remark 2

lemma *unique-transversal-segment-intersection:*
assumes *transversal-segment a b*
assumes $\{0..\} \subseteq \text{existence-ivl0 } x$
assumes $u \in \omega\text{-limit-set } x \cap \{a <--< b\}$
shows $\omega\text{-limit-set } x \cap \{a <--< b\} = \{u\}$
 <proof>

Adapted from Perko Section 3.7 Lemma 4 (+ Chicone)

lemma *periodic-imp- ω -limit-set:*
assumes *compact K K \subseteq X*
assumes $x \in X$ *trapped-forward x K*
assumes *periodic-orbit y*
 $\text{flow0 } y \text{ 'UNIV} \subseteq \omega\text{-limit-set } x$
shows $\text{flow0 } y \text{ 'UNIV} = \omega\text{-limit-set } x$

<proof>

end context *c1-on-open-R2* **begin**

lemma *α -limit-crossings:*

assumes *transversal-segment* $a\ b$
assumes *pos-ex:* $\{..0\} \subseteq \textit{existence-ivl0}\ x$
assumes *α -limit-point* $x\ p$
assumes $p \in \{a < - - < b\}$
obtains s **where**
 $s \longrightarrow -\infty$
 $(\textit{flow0}\ x \circ s) \longrightarrow p$
 $\forall_F n$ *in sequentially.*
 $\textit{flow0}\ x\ (s\ n) \in \{a < - - < b\} \wedge$
 $s\ n \in \textit{existence-ivl0}\ x$

<proof>

If a positive limit point has a regular point in its positive limit set then it is periodic

lemma *ω -limit-point- ω -limit-set-regular-imp-periodic:*

assumes *compact* $K\ K \subseteq X$
assumes $x \in X$ *trapped-forward* $x\ K$
assumes $y: y \in \omega\text{-limit-set}\ x\ f\ y \neq 0$
assumes $z: z \in \omega\text{-limit-set}\ y \cup \alpha\text{-limit-set}\ y\ f\ z \neq 0$
shows *periodic-orbit* $y \wedge \textit{flow0}\ y\ 'UNIV = \omega\text{-limit-set}\ x$

<proof>

6.8 Poincare Bendixson Theorems

Perko Section 3.7 Theorem 1

theorem *poincare-bendixson:*

assumes *compact* $K\ K \subseteq X$
assumes $x \in X$ *trapped-forward* $x\ K$
assumes $0 \notin f\ '(\omega\text{-limit-set}\ x)$
obtains y **where** *periodic-orbit* y
 $\textit{flow0}\ y\ 'UNIV = \omega\text{-limit-set}\ x$

<proof>

lemma *fixed-point-in- ω -limit-set-imp- ω -limit-set-singleton-fixed-point:*

assumes *compact* $K\ K \subseteq X$
assumes $x \in X$ *trapped-forward* $x\ K$
assumes $fp: yfp \in \omega\text{-limit-set}\ x\ f\ yfp = 0$
assumes $zpx: z \in \omega\text{-limit-set}\ x$
assumes *finite-fp:* $\textit{finite}\ \{y \in K. f\ y = 0\}$ (**is finite** ? S)
shows $(\exists p1 \in \omega\text{-limit-set}\ x. f\ p1 = 0 \wedge \omega\text{-limit-set}\ z = \{p1\}) \wedge$
 $(\exists p2 \in \omega\text{-limit-set}\ x. f\ p2 = 0 \wedge \alpha\text{-limit-set}\ z = \{p2\})$

<proof>

end context *c1-on-open-R2* **begin**

Perko Section 3.7 Theorem 2

theorem *poincare-bendixson-general*:

assumes *compact* $K \subseteq X$

assumes $x \in X$ *trapped-forward* $x \in K$

assumes $S = \{y \in K. f \cdot y = 0\}$ *finite* S

shows

$(\exists y \in S. \omega\text{-limit-set } x = \{y\}) \vee$

$(\exists y. \text{periodic-orbit } y \wedge$

$\text{flow0 } y \text{ ' UNIV } = \omega\text{-limit-set } x) \vee$

$(\exists P R. \omega\text{-limit-set } x = P \cup R \wedge$

$P \subseteq S \wedge 0 \notin f \text{ ' } R \wedge R \neq \{\}) \wedge$

$(\forall z \in R.$

$(\exists p1 \in P. \omega\text{-limit-set } z = \{p1\}) \wedge$

$(\exists p2 \in P. \alpha\text{-limit-set } z = \{p2\}))$)

<proof>

corollary *poincare-bendixson-applied*:

assumes *compact* $K \subseteq X$

assumes $K \neq \{\}$ *positively-invariant* K

assumes $0 \notin f \text{ ' } K$

obtains y **where** *periodic-orbit* $y \text{ flow0 } y \text{ ' UNIV } \subseteq K$

<proof>

definition *limit-cycle* $y \longleftrightarrow$

periodic-orbit $y \wedge$

$(\exists x. x \notin \text{flow0 } y \text{ ' UNIV} \wedge$

$(\text{flow0 } y \text{ ' UNIV} = \omega\text{-limit-set } x \vee \text{flow0 } y \text{ ' UNIV} = \alpha\text{-limit-set } x))$

corollary *poincare-bendixson-limit-cycle*:

assumes *compact* $K \subseteq X$

assumes $x \in K$ *positively-invariant* K

assumes $0 \notin f \text{ ' } K$

assumes *rev.flow0* $x \notin K$

obtains y **where** *limit-cycle* $y \text{ flow0 } y \text{ ' UNIV } \subseteq K$

<proof>

end

end

theory *Affine-Arithmetic-Misc*

imports *HOL-ODE-Numerics.ODE-Numerics*

begin

7 Branch-And-Bound Arithmetic

primrec *prove-nonneg*::(nat * nat * string) list \Rightarrow nat \Rightarrow nat \Rightarrow slp \Rightarrow real aform list list \Rightarrow bool **where**

prove-nonneg prnt 0 p slp X = (let - = if prnt \neq [] then print (STR "# depth limit exceeded"') else () in False)

| *prove-nonneg* prnt (Suc i) p slp XXS =
 (case XXS of [] \Rightarrow True | (X#XS) \Rightarrow
 let RS = approx-slp-outer p 1 slp X
 in if RS \neq None \wedge Inf-aform' p (hd (the RS)) \geq 0
 then
 let - = if prnt \neq [] then print (STR "# Success"') else ();
 - = if prnt \neq [] then print (String.implode ((shows "# " o shows-box-of-aforms-hr X) ')) else ();
 - = fold ($\lambda(a, b, c) \cdot$ print (String.implode (shows-segments-of-aform a b X c '))) prnt ()
 in *prove-nonneg* prnt i p slp XS
 else let - = if prnt \neq [] then print (STR "# Split"') else () in case split-aforms-largest-uncond X of (a, b) \Rightarrow
prove-nonneg prnt i p slp (a#b#XS))

lemma *prove-nonneg-simps*[simp]:
prove-nonneg prnt 0 p slp X = False
prove-nonneg prnt (Suc i) p slp XXS =
 (case XXS of [] \Rightarrow True | (X#XS) \Rightarrow
 let RS = approx-slp-outer p 1 slp X
 in if RS \neq None \wedge Inf-aform' p (hd (the RS)) \geq 0
 then *prove-nonneg* prnt i p slp XS
 else case split-aforms-largest-uncond X of (a, b) \Rightarrow *prove-nonneg* prnt i p slp (a#b#XS))
 <proof>

lemmas [simp del] = *prove-nonneg.simps*

lemma *split-aforms-lemma*:
fixes xs::real list
assumes split-aforms XS i = (YS, ZS)
assumes xs \in Joints XS
shows xs \in Joints YS \cup Joints ZS
 <proof>

lemma *prove-nonneg-empty*[simp]: *prove-nonneg* prnt (Suc i) p slp []
 <proof>

lemma *prove-nonneg-fuel-mono*:
prove-nonneg prnt (Suc i) p (slp-of-fas [fa]) YSS
 if *prove-nonneg* prnt i p (slp-of-fas [fa]) YSS
 <proof>

lemma *prove-nonneg-mono*:
prove-nonneg prnt i p (slp-of-fas [fa]) YSS if prove-nonneg prnt i p (slp-of-fas [fa]) (YS # YSS)
<proof>

lemma *prove-nonneg*:
assumes *prove-nonneg prnt i p (slp-of-fas [fa]) XSS*
shows $\forall XS \in \text{set } XSS. \forall xs \in \text{Joints } XS. \text{interpret-floatarith } fa \text{ } xs \geq 0$
<proof>

end

8 Examples

theory *Examples*
imports *Poincare-Bendixson*
HOL-ODE-Numerics.ODE-Numerics
Affine-Arithmetic-Misc
begin

8.1 Simple

context
begin

coordinate functions

definition $cx \ x \ y = -y + x * (1 - x^2 - y^2)$

definition $cy \ x \ y = x + y * (1 - x^2 - y^2)$

lemmas *c-defs = cx-def cy-def*

partial derivatives

definition $C11 :: \text{real} \Rightarrow \text{real} \Rightarrow \text{real}$ **where** $C11 \ x \ y = 1 - 3 * x^2 - y^2$

definition $C12 :: \text{real} \Rightarrow \text{real} \Rightarrow \text{real}$ **where** $C12 \ x \ y = -1 - 2 * x * y$

definition $C21 :: \text{real} \Rightarrow \text{real} \Rightarrow \text{real}$ **where** $C21 \ x \ y = 1 - 2 * x * y$

definition $C22 :: \text{real} \Rightarrow \text{real} \Rightarrow \text{real}$ **where** $C22 \ x \ y = 1 - x^2 - 3 * y^2$

lemmas *C-partials = C11-def C12-def C21-def C22-def*

Jacobian as linear map

definition $C :: \text{real} \Rightarrow \text{real} \Rightarrow (\text{real} \times \text{real}) \Rightarrow_L (\text{real} \times \text{real})$ **where**

$C \ x \ y = \text{blinfun-of-matrix}$

$((\lambda-. 0)$

$((1,0) := (\lambda-. 0)((1, 0) := C11 \ x \ y, (0, 1) := C12 \ x \ y),$

$(0, 1) := (\lambda-. 0)((1, 0) := C21 \ x \ y, (0, 1) := C22 \ x \ y)))$

lemma *C-simp[simp]*: $\text{blinfun-apply } (C \ x \ y) \ (dx, dy) =$
 $(dx * C11 \ x \ y + dy * C12 \ x \ y,$

$dx * C21 x y + dy * C22 x y$
 $\langle proof \rangle$

lemma *C-continuous*[*continuous-intros*]:
continuous-on S ($\lambda x. local.C (f x) (g x)$)
if *continuous-on S f continuous-on S g*
 $\langle proof \rangle$

interpretation *c: c1-on-open-R2* $\lambda(x::real, y::real). (cx x y, cy x y)::real*real$
 $\lambda(x, y). C x y UNIV$
 $\langle proof \rangle$

definition *trapC* = *cball* ($0::real, 0::real$) 2 - *ball* ($0::real, 0::real$) (1/2)

lemma *trapC-eq*:
shows $trapC = \{p. (fst p)^2 + (snd p)^2 - 4 \leq 0\} \cap \{p. 1/4 - ((fst p)^2 + (snd p)^2) \leq 0\}$
 $\langle proof \rangle$

lemma *x-in-trapC*:
shows $(2, 0) \in trapC$
 $\langle proof \rangle$

lemma *compact-trapC*:
shows *compact trapC*
 $\langle proof \rangle$

lemma *nonempty-trapC*:
shows $trapC \neq \{\}$
 $\langle proof \rangle$

lemma *origin-fixpoint*:
assumes $(\lambda(x, y). (cx x y, cy x y)) (a, b) = 0$
shows $a = (0::real) b = (0::real)$
 $\langle proof \rangle$

lemma *origin-not-trapC*:
shows $0 \notin trapC$
 $\langle proof \rangle$

lemma *regular-trapC*:
shows $0 \notin (\lambda(x, y). (cx x y, cy x y)) ` trapC$
 $\langle proof \rangle$

lemma *positively-invariant-outer*:
shows *c.positively-invariant* $\{p. (\lambda p. (fst p)^2 + (snd p)^2 - 4) p \leq 0\}$
 $\langle proof \rangle$

lemma *positively-invariant-inner*:

shows *c.positively-invariant* $\{p. (\lambda p. 1/4 - ((fst\ p)^2 + (snd\ p)^2))\ p \leq 0\}$
<proof>

lemma *positively-invariant-trapC*:

shows *c.positively-invariant trapC*
<proof>

theorem *c-has-periodic-orbit*:

obtains *y where c.periodic-orbit y c.flow0 y ' UNIV \subseteq trapC*
<proof>

Real-Arithmetic

schematic-goal *c-fas*:

$[-(-(X!1) + (X!0) * (1 - (X!0)^2 - (X!1)^2)), -((X!0) + (X!1) * (1 - (X!0)^2 - (X!1)^2))] = interpret-floatariths\ ?fas\ X$
<proof>

concrete-definition *c-fas uses c-fas*

interpretation *crev: ode-interpretation true-form UNIV c-fas*

$-(\lambda(x, y). (cx\ x\ y, cy\ x\ y)::real*real)$
d::2 for d
<proof>

lemma *crev: $t \in \{1/8 .. 1/8\} \longrightarrow (x, y) \in \{(2, 0) .. (2, 0)\} \longrightarrow$*

$t \in c.rev.existence-ivl0\ (x, y) \wedge c.rev.flow0\ (x, y)\ t \in \{(5.15, -0.651)..(5.18, -0.647)\}$
<proof>

theorem *c-has-limit-cycle*:

obtains *y where c.limit-cycle y range (c.flow0 y) \subseteq trapC*
<proof>

end

8.2 Glycolysis

Strogatz, Example 7.3.2

context

begin

coordinate functions

definition *gx* $x\ y = -x + 0.08 * y + x^2 * y$

definition *gy* $x\ y = 0.6 - 0.08 * y - x^2 * y$

lemmas *g-defs = gx-def gy-def*

partial derivatives

definition $A11::real \Rightarrow real \Rightarrow real$ **where** $A11\ x\ y = -1 + 2 * x * y$

definition $A12::real \Rightarrow real \Rightarrow real$ **where** $A12\ x\ y = (0.08 + x^2)$

definition $A21::real \Rightarrow real \Rightarrow real$ **where** $A21\ x\ y = -2*x*y$

definition $A22::real \Rightarrow real \Rightarrow real$ **where** $A22\ x\ y = -(0.08 + x^2)$

lemmas A -partials = $A11$ -def $A12$ -def $A21$ -def $A22$ -def

Jacobian as linear map

definition $A :: real \Rightarrow real \Rightarrow (real \times real) \Rightarrow_L (real \times real)$ **where**

$A\ x\ y = blinfun-of-matrix$

$((\lambda-. 0)$

$((1, 0) := (\lambda-. 0)((1, 0) := A11\ x\ y, (0, 1) := A12\ x\ y),$

$(0, 1) := (\lambda-. 0)((1, 0) := A21\ x\ y, (0, 1) := A22\ x\ y))$

lemma A -simp[simp]: $blinfun-apply\ (A\ x\ y)\ (dx, dy) =$

$(dx * A11\ x\ y + dy * A12\ x\ y,$

$dx * A21\ x\ y + dy * A22\ x\ y)$

$\langle proof \rangle$

lemma A -continuous[continuous-intros]:

$continuous-on\ S\ (\lambda x. local.A\ (f\ x)\ (g\ x))$

if $continuous-on\ S\ f\ continuous-on\ S\ g$

$\langle proof \rangle$

interpretation g : $c1$ -on-open- $R^2\ \lambda(x::real, y::real). (gx\ x\ y, gy\ x\ y)::real*real$

$\lambda(x, y). A\ x\ y\ UNIV$

$\langle proof \rangle$

definition $(pos\ quad::(real \times real)\ set) = \{p . -\ snd\ p \leq 0\} \cap \{p . -\ fst\ p \leq 0\}$

definition $(trapG1::(real \times real)\ set) = pos\ quad \cap (\{p. (snd\ p) - 751/100 \leq 0\}$
 $\cap \{p. (fst\ p) + (snd\ p) - 812/100 \leq 0\})$

lemma $positively$ -invariant- y :

shows g . $positively$ -invariant $\{p . -\ snd\ p \leq 0\}$

$\langle proof \rangle$

lemma $positively$ -invariant- pos -quad:

shows g . $positively$ -invariant $pos\ quad$

$\langle proof \rangle$

lemma $positively$ -invariant- y -upper:

shows g . $positively$ -invariant $\{p. (snd\ p) - 751/100 \leq 0\}$

$\langle proof \rangle$

lemma $arith2$:

shows $(y::real) \leq 751/100 \wedge x + (y::real) = 812/100 \implies 3/5 - (x::real) < 0$

$\langle proof \rangle$

lemma *positively-invariant-trapG1*:
shows *g.positively-invariant trapG1*
 ⟨*proof*⟩

definition *p1* (*x::real*) (*y::real*) = $-(21/34) - (69*x)/38 + (19*x^2)/15 - (9*x^3)/28 - (6*x^4)/43 + (14*y)/29 + (31*x*y)/21 + (182*x^2*y)/47 - (35*x^3*y)/16 - (3*y^2)/17 - (2*x*y^2)/9 - (31*x^2*y^2)/20 + y^3/102 + (x*y^3)/59$

definition *p1d* $x\ xa = 38 * (fst\ xa * fst\ x) / 15 - 69 * fst\ xa / 38 - 27 * (fst\ xa * (fst\ x)^2) / 28 - 24 * (fst\ xa * fst\ x^3) / 43 + 14 * snd\ xa / 29 + (651 * (fst\ x * snd\ xa) + 651 * (fst\ xa * snd\ x)) / 441 + (8554 * ((fst\ x)^2 * snd\ xa) + 17108 * (fst\ xa * (fst\ x * snd\ x))) / 2209 - (560 * (fst\ x^3 * snd\ xa) + 1680 * (fst\ xa * ((fst\ x)^2 * snd\ x))) / 256 - 6 * (snd\ xa * snd\ x) / 17 - (36 * (fst\ x * (snd\ xa * snd\ x)) + 18 * (fst\ xa * (snd\ x)^2)) / 81 - (1240 * ((fst\ x)^2 * (snd\ xa * snd\ x)) + 1240 * (fst\ xa * (fst\ x * (snd\ x)^2))) / 400 + snd\ xa * (snd\ x)^2 / 34 + (177 * (fst\ x * (snd\ xa * (snd\ x)^2)) + fst\ xa * snd\ x^3 * 59) / 3481$

lemma *p1-has-derivative*:
shows $((\lambda x. p1\ (fst\ x)\ (snd\ x))\ has\ derivative\ p1d\ x)\ (at\ x)$
 ⟨*proof*⟩

lemma *p1-not-equil*:
shows $p1\ x\ y \leq 0 \implies gx\ x\ y \neq 0 \vee gy\ x\ y \neq 0$
 ⟨*proof*⟩

definition *trapG* = $trapG1 \cap \{p. p1\ (fst\ p)\ (snd\ p) \leq 0\}$

Real-Arithmetic

definition *g-arith* $a\ b = (- (27 / 25) - a^2 + 2 * a * b) * p1\ a\ b - p1d\ (a, b)$

($gx\ a\ b$, $gy\ a\ b$)

schematic-goal *g-arith-fas*:

[*g-arith* ($X!0$) ($X!1$)] = *interpret-floatariths ?fas X*
<proof>

concrete-definition *g-arith-fas uses g-arith-fas*

lemma *list-interval2*: *list-interval* [a , b] [c , d] = $\{[x, y] \mid x\ y.\ x \in \{a \dots c\} \wedge y \in \{b \dots d\}\}$
<proof>

lemma *g-arith-nonneg*: *g-arith* $a\ b \geq 0$
if $a: 0 \leq a\ a \leq 8.24$ and $b: 0 \leq b\ b \leq 7.51$
<proof>

lemma *trap-arithmetic*:
 $p1d\ (a, b)\ (gx\ a\ b, gy\ a\ b) \leq (- (27 / 25) - a^2 + 2 * a * b) * p1\ a\ b$ if (a, b)
 $\in\ trapG1$
<proof>

lemma *positively-invariant-trapG*:
shows *g.positively-invariant trapG*
<proof>

lemma *regular-trapG*:
shows $0 \notin (\lambda(x, y). (gx\ x\ y, gy\ x\ y))\ ` trapG$
<proof>

lemma *arith*:
 $\bigwedge a\ b::real. 0 \leq b \implies$
 $0 \leq a \implies$
 $b * 100 \leq 751 \implies$
 $a * 25 + b * 25 \leq 203 \implies norm\ a + norm\ b \leq 20$
<proof>

lemma *trapG1-subset*:
shows $trapG1 \subseteq cball\ (0::real \times real)\ 20$
<proof>

lemma *compact-subset-closed*:
assumes *compact S closed T*
assumes $T \subseteq S$
shows *compact T*
<proof>

lemma *compact-trapG1*:
shows *compact trapG1*
<proof>

lemma *compact-trapG*:
shows *compact trapG*
 ⟨*proof*⟩

lemma *x-in-trapG*:
shows $(1,0) \in \text{trap}G$
 ⟨*proof*⟩

schematic-goal *g-fas*:

$$[-(- (X!0) + 8 / 100 * (X!1) + (X!0)^2 * (X!1)), -(6 / 10 - 8 / 100 * (X!1) - (X!0)^2 * (X!1))] = \text{interpret-floatariths } ?\text{fas } X$$

 ⟨*proof*⟩

concrete-definition *g-fas uses g-fas*

interpretation *grev: ode-interpretation true-form UNIV g-fas*
 $-(\lambda(x, y). (gx\ x\ y, gy\ x\ y)::\text{real}*\text{real})$
d::2 for d
 ⟨*proof*⟩

lemma *grev: $t \in \{1/8 .. 1/8\} \longrightarrow (x, y) \in \{(1, 0) .. (1, 0)\} \longrightarrow t \in g.\text{rev}.\text{existence-ivl0 } (x, y) \wedge g.\text{rev}.\text{flow0 } (x, y) t \in \{(1.1, -0.09) .. (1.2, -0.08)\}$*
 ⟨*proof*⟩

theorem *g-has-limit-cycle*:
obtains *y where* $g.\text{limit-cycle } y \text{ range } (g.\text{flow0 } y) \subseteq \text{trap}G$
 ⟨*proof*⟩

end

end