

Pick's Theorem

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Abstract

We formalize Pick's theorem for finding the area of a simple polygon whose vertices are integral lattice points [1]. We are inspired by John Harrison's formalization of Pick's theorem in HOL Light [2], but tailor our proof approach to avoid a primary challenge point in his formalization, which is proving that any polygon with more than three vertices can be split (in its interior) by a line between some two vertices. Our formalization involves augmenting the existing geometry libraries in various foundational ways (e.g., by adding the definition of a polygon and formalizing some key properties thereof).

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theory <i>Integral-Matrix</i>	
imports	
<i>Complex-Main</i>	
<i>HOL-Analysis.Finite-Cartesian-Product</i>	
<i>HOL-Analysis.Linear-Algebra</i>	
<i>HOL-Analysis.Determinants</i>	
begin	

1 Misc. Linear Algebra Setup

lemma *vec-scaleR-2*: $(c::real) *_R ((vector [a, b])::real^2) = vector [a * c, b * c]$
<proof>

definition *is-int* :: $real \Rightarrow bool$ **where**
 $is-int\ x \longleftrightarrow (\exists n::int. x = n)$

lemma *is-int-sum*: $is-int\ x \wedge is-int\ y \longrightarrow is-int\ (x + y)$
<proof>

lemma *is-int-minus*: $is-int\ x \wedge is-int\ y \longrightarrow is-int\ (x - y)$
<proof>

lemma *is-int-mult*: $is-int\ x \wedge is-int\ y \longrightarrow is-int\ (x * y)$
<proof>

definition *integral-vec* :: $real^2 \Rightarrow bool$ **where**
 $integral-vec\ v \longleftrightarrow (is-int\ (v\$1) \wedge is-int\ (v\$2))$

lemma *integral-vec-sum*: $integral-vec\ v \wedge integral-vec\ w \longrightarrow integral-vec\ (v + w)$
<proof>

lemma *integral-vec-minus*: $integral-vec\ v \longrightarrow integral-vec\ (-v)$
<proof>

lemma *real-2-inner*:
shows $((vector [a, b])::(real^2)) \cdot ((vector [c, d])::(real^2)) = a*c + b*d$
(is $?v \cdot ?w = a*c + b*d$
<proof>

lemma *integral-vec-2*:
fixes $a\ b :: int$
assumes $v = vector [a, b]$
shows $integral-vec\ v$
<proof>

definition *matrix-inv* :: $real^2 \Rightarrow real^2 \Rightarrow bool$ **where**
 $matrix-inv\ A\ A' \longleftrightarrow (A ** A' = mat\ 1 \wedge A' ** A = mat\ 1)$

lemma *mat-vec-mult-2*:
fixes $v :: real^2$ **and**
 $T :: real^2 \Rightarrow real^2$
defines $x: x \equiv v\$1$ **and** $y: y \equiv v\$2$ **and**
 $a: a \equiv T\$1\1 **and** $b: b \equiv T\$1\2 **and**
 $c: c \equiv T\$2\1 **and** $d: d \equiv T\$2\2
shows $(T *v v) = vector [x*a + y*b, x*c + y*d]$
<proof>

definition *integral-mat* :: $\text{real}^{\mathbb{Q}^2} \Rightarrow \text{bool}$ **where**
integral-mat $T \iff (\forall v. \text{integral-vec } v \longrightarrow \text{integral-vec } (T * v))$

definition *integral-mat-surj* :: $\text{real}^{\mathbb{Q}^2} \Rightarrow \text{bool}$ **where**
integral-mat-surj $T \iff (\forall v. \text{integral-vec } v \longrightarrow (\exists w. \text{integral-vec } w \wedge T * v = w))$

definition *integral-mat-bij* :: $\text{real}^{\mathbb{Q}^2} \Rightarrow \text{bool}$ **where**
integral-mat-bij $T \iff \text{integral-mat } T \wedge \text{integral-mat-surj } T$

lemma *integral-mat-integral-vec*: $\text{integral-mat } A \longrightarrow \text{integral-vec } v \longrightarrow \text{integral-vec } (A * v)$
 ⟨*proof*⟩

lemma *integral-mat-int-entries*:
fixes $T :: \text{real}^{\mathbb{Q}^2}$
assumes *integral-mat* T
defines $a: a \equiv T\$1\1 **and** $b: b \equiv T\$1\2 **and**
 $c: c \equiv T\$2\1 **and** $d: d \equiv T\$2\2
shows $\text{is-int } a \wedge \text{is-int } b \wedge \text{is-int } c \wedge \text{is-int } d$
 ⟨*proof*⟩

2 Integral Bijective Matrix Determinant

lemma *integral-mat-int-det*:
fixes $T :: \text{real}^{\mathbb{Q}^2}$
assumes *integral-mat* T
shows $\text{is-int } (\text{det } T)$
 ⟨*proof*⟩

lemma *integral-mat-bij-inv*:
fixes $T :: \text{real}^{\mathbb{Q}^2}$
assumes *integral-mat-bij* T
obtains T_{inv} **where** $\text{invertible } T \wedge \text{integral-mat-bij } T_{\text{inv}} \wedge \text{matrix-inv } T T_{\text{inv}}$
 ⟨*proof*⟩

lemma *integral-mat-bij-det-pm1*:
fixes $T :: \text{real}^{\mathbb{Q}^2}$
assumes *integral-mat-bij* T
shows $\text{det } T = 1 \vee \text{det } T = -1$
 ⟨*proof*⟩

end
theory *Polygon-Jordan-Curve*
imports
HOL-Analysis.Cartesian-Space
HOL-Analysis.Path-Connected

begin

3 Polygon Definitions

type-synonym $R\text{-to-}R2 = (real \Rightarrow real^2)$

definition $closed\text{-}path :: R\text{-to-}R2 \Rightarrow bool$ **where**
 $closed\text{-}path\ g \longleftrightarrow path\ g \wedge pathstart\ g = pathfinish\ g$

definition $path\text{-}inside :: R\text{-to-}R2 \Rightarrow (real^2)$ *set* **where**
 $path\text{-}inside\ g = inside\ (path\text{-}image\ g)$

definition $path\text{-}outside :: R\text{-to-}R2 \Rightarrow (real^2)$ *set* **where**
 $path\text{-}outside\ g = outside\ (path\text{-}image\ g)$

fun $make\text{-}polygonal\text{-}path :: (real^2)$ *list* $\Rightarrow R\text{-to-}R2$ **where**
 $make\text{-}polygonal\text{-}path\ [] = linepath\ 0\ 0$
 $| make\text{-}polygonal\text{-}path\ [a] = linepath\ a\ a$
 $| make\text{-}polygonal\text{-}path\ [a,b] = linepath\ a\ b$
 $| make\text{-}polygonal\text{-}path\ (a \# b \# xs) = (linepath\ a\ b) +++ make\text{-}polygonal\text{-}path\ (b \# xs)$

definition $polygonal\text{-}path :: R\text{-to-}R2 \Rightarrow bool$ **where**
 $polygonal\text{-}path\ g \longleftrightarrow g \in make\text{-}polygonal\text{-}path\ \{xs :: (real^2)$ *list*. $True\}$

definition $all\text{-}integral :: (real^2)$ *list* $\Rightarrow bool$ **where**
 $all\text{-}integral\ l = (\forall x \in set\ l. integral\text{-}vec\ x)$

definition $polygon :: R\text{-to-}R2 \Rightarrow bool$ **where**
 $polygon\ g \longleftrightarrow polygonal\text{-}path\ g \wedge simple\text{-}path\ g \wedge closed\text{-}path\ g$

definition $integral\text{-}polygon :: R\text{-to-}R2 \Rightarrow bool$ **where**
 $integral\text{-}polygon\ g \longleftrightarrow$
 $(polygon\ g \wedge (\exists vts. g = make\text{-}polygonal\text{-}path\ vts \wedge all\text{-}integral\ vts))$

definition $make\text{-}triangle :: real^2 \Rightarrow real^2 \Rightarrow real^2 \Rightarrow R\text{-to-}R2$ **where**
 $make\text{-}triangle\ a\ b\ c = make\text{-}polygonal\text{-}path\ [a, b, c, a]$

definition $polygon\text{-}of :: R\text{-to-}R2 \Rightarrow (real^2)$ *list* $\Rightarrow bool$ **where**
 $polygon\text{-}of\ p\ vts \longleftrightarrow polygon\ p \wedge p = make\text{-}polygonal\text{-}path\ vts$

definition $good\text{-}linepath :: real^2 \Rightarrow real^2 \Rightarrow (real^2)$ *list* $\Rightarrow bool$ **where**
 $good\text{-}linepath\ a\ b\ vts \longleftrightarrow (let\ p = make\text{-}polygonal\text{-}path\ vts\ in$
 $a \neq b \wedge \{a, b\} \subseteq set\ vts \wedge path\text{-}image\ (linepath\ a\ b) \subseteq path\text{-}inside\ p \cup \{a, b\})$

definition $good\text{-}polygonal\text{-}path :: real^2 \Rightarrow (real^2)$ *list* $\Rightarrow real^2 \Rightarrow (real^2)$ *list*

\Rightarrow **bool where**
good-polygonal-path a *cutvts* b *vts* \longleftrightarrow (
 let $p = \text{make-polygonal-path } vts$ in
 let $p\text{-cut} = \text{make-polygonal-path } ([a] @ \text{cutvts} @ [b])$ in
 ($a \neq b \wedge \{a, b\} \subseteq \text{set } vts \wedge \text{path-image } (p\text{-cut}) \subseteq \text{path-inside } p \cup \{a, b\} \wedge$
loop-free $p\text{-cut}$)

4 Jordan Curve Theorem for Polygons

definition *inside-outside* :: $R\text{-to-}R^2 \Rightarrow (\text{real}^2)$ *set* $\Rightarrow (\text{real}^2)$ *set* \Rightarrow **bool where**
inside-outside p *ins* *outs* \longleftrightarrow
 ($ins \neq \{\}$ \wedge *open* $ins \wedge$ *connected* $ins \wedge$
 $outs \neq \{\}$ \wedge *open* $outs \wedge$ *connected* $outs \wedge$
bounded $ins \wedge \neg$ *bounded* $outs \wedge$
 $ins \cap outs = \{\} \wedge ins \cup outs = - \text{path-image } p \wedge$
frontier $ins = \text{path-image } p \wedge$ *frontier* $outs = \text{path-image } p$)

lemma *Jordan-inside-outside-real2*:

fixes $p :: \text{real} \Rightarrow \text{real}^2$

assumes *simple-path* p *pathfinish* $p = \text{pathstart } p$

shows $inside(\text{path-image } p) \neq \{\} \wedge$
 $open(inside(\text{path-image } p)) \wedge$
 $connected(inside(\text{path-image } p)) \wedge$
 $outside(\text{path-image } p) \neq \{\} \wedge$
 $open(outside(\text{path-image } p)) \wedge$
 $connected(outside(\text{path-image } p)) \wedge$
 $bounded(inside(\text{path-image } p)) \wedge$
 $\neg bounded(outside(\text{path-image } p)) \wedge$
 $inside(\text{path-image } p) \cap outside(\text{path-image } p) = \{\} \wedge$
 $inside(\text{path-image } p) \cup outside(\text{path-image } p) =$
 $- \text{path-image } p \wedge$
 $frontier(inside(\text{path-image } p)) = \text{path-image } p \wedge$
 $frontier(outside(\text{path-image } p)) = \text{path-image } p$

<proof>

lemma *inside-outside-polygon*:

fixes $p :: R\text{-to-}R^2$

assumes *polygon*: *polygon* p

shows *inside-outside* p (*path-inside* p) (*path-outside* p)

<proof>

lemma *inside-outside-unique*:

fixes $p :: R\text{-to-}R^2$

assumes *polygon* p

assumes *io1*: *inside-outside* p *inside1* *outside1*

assumes *io2*: *inside-outside* p *inside2* *outside2*

shows $inside1 = inside2 \wedge outside1 = outside2$

<proof>

lemma *polygon-jordan-curve*:

fixes $p :: R\text{-to-}R^2$

assumes *polygon* p

obtains *inside outside* **where**

inside-outside p *inside outside*

<proof>

lemma *connected-component-image*:

fixes $f :: 'a::\text{euclidean-space} \Rightarrow 'b::\text{euclidean-space}$

assumes *linear* f *bij* f

shows $f^{-1}(\text{connected-component-set } S \ x) = \text{connected-component-set } (f^{-1} \ S) \ (f \ x)$

<proof>

lemma *bounded-map*:

fixes $f :: 'a::\text{euclidean-space} \Rightarrow 'b::\text{euclidean-space}$

assumes *linear* f *bij* f

shows *bounded* $(f^{-1} \ S) = \text{bounded } S$

<proof>

lemma *inside-bijective-linear-image*:

fixes $f :: 'a::\text{euclidean-space} \Rightarrow 'b::\text{euclidean-space}$

fixes $c :: \text{real} \Rightarrow 'a$

assumes *c-simple: path* c

assumes *linear* f *bij* f

shows *inside* $(f^{-1}(\text{path-image } c)) = f^{-1}(\text{inside}(\text{path-image } c))$

<proof>

lemma *bij-image-intersection*:

assumes *path-image* $c1 \cap \text{path-image } c2 = S$

assumes *bij* f

assumes $c \in \text{path-image } (f \circ c1) \cap \text{path-image } (f \circ c2)$

shows $c \in f^{-1} \ S$

<proof>

theorem (*in* *c1-on-open-R2*) *split-inside-simple-closed-curve-locale*:

fixes $c :: \text{real} \Rightarrow 'a$

assumes *c1-simple: simple-path* $c1$ **and** *c1-start: pathstart* $c1 = a$ **and** *c1-end: pathfinish* $c1 = b$

assumes *c2-simple: simple-path* $c2$ **and** *c2-start: pathstart* $c2 = a$ **and** *c2-end: pathfinish* $c2 = b$

assumes *c-simple: simple-path* c **and** *c-start: pathstart* $c = a$ **and** *c-end: pathfinish* $c = b$

assumes *a-neq-b: a* $\neq b$

and *c1c2: path-image* $c1 \cap \text{path-image } c2 = \{a, b\}$

and *c1c: path-image* $c1 \cap \text{path-image } c = \{a, b\}$

and $c2c$: $path\text{-}image\ c2 \cap path\text{-}image\ c = \{a,b\}$
and $ne\text{-}12$: $path\text{-}image\ c \cap inside(path\text{-}image\ c1 \cup path\text{-}image\ c2) \neq \{\}$
obtains $inside(path\text{-}image\ c1 \cup path\text{-}image\ c) \cap inside(path\text{-}image\ c2 \cup path\text{-}image\ c) = \{\}$
 $inside(path\text{-}image\ c1 \cup path\text{-}image\ c) \cup inside(path\text{-}image\ c2 \cup path\text{-}image\ c) \cup$
 $(path\text{-}image\ c - \{a,b\}) = inside(path\text{-}image\ c1 \cup path\text{-}image\ c2)$
 $\langle proof \rangle$

lemma *split-inside-simple-closed-curve-real2*:

fixes $c :: real \Rightarrow real^2$
assumes $c1\text{-}simple$: $simple\text{-}path\ c1$ **and** $c1\text{-}start$: $pathstart\ c1 = a$ **and** $c1\text{-}end$:
 $pathfinish\ c1 = b$
assumes $c2\text{-}simple$: $simple\text{-}path\ c2$ **and** $c2\text{-}start$: $pathstart\ c2 = a$ **and** $c2\text{-}end$:
 $pathfinish\ c2 = b$
assumes $c\text{-}simple$: $simple\text{-}path\ c$ **and** $c\text{-}start$: $pathstart\ c = a$ **and** $c\text{-}end$: $pathfinish\ c = b$
assumes $a\text{-}neq\text{-}b$: $a \neq b$
and $c1c2$: $path\text{-}image\ c1 \cap path\text{-}image\ c2 = \{a,b\}$
and $c1c$: $path\text{-}image\ c1 \cap path\text{-}image\ c = \{a,b\}$
and $c2c$: $path\text{-}image\ c2 \cap path\text{-}image\ c = \{a,b\}$
and $ne\text{-}12$: $path\text{-}image\ c \cap inside(path\text{-}image\ c1 \cup path\text{-}image\ c2) \neq \{\}$
obtains $inside(path\text{-}image\ c1 \cup path\text{-}image\ c) \cap inside(path\text{-}image\ c2 \cup path\text{-}image\ c) = \{\}$
 $inside(path\text{-}image\ c1 \cup path\text{-}image\ c) \cup inside(path\text{-}image\ c2 \cup path\text{-}image\ c) \cup$
 $(path\text{-}image\ c - \{a,b\}) = inside(path\text{-}image\ c1 \cup path\text{-}image\ c2)$
 $\langle proof \rangle$

end

theory *Polygon-Lemmas*

imports

Polygon-Jordan-Curve
HOL-Library.Sublist
HOL.Set-Interval
HOL.Fun

begin

5 Properties of make polygonal path, pathstart and pathfinish of a polygon

lemma *make-polygonal-path-induct*[*case-names Empty Single Two Multiple*]:

fixes $ell :: (real^2) list$
assumes $empty$: $\bigwedge ell. ell = [] \implies P\ ell$
and $single$: $\bigwedge ell. \llbracket length\ ell = 1 \rrbracket \implies P\ ell$
and two : $\bigwedge ell. \llbracket length\ ell = 2 \rrbracket \implies P\ ell$
and $multiple$: $\bigwedge ell.$

$\llbracket \text{length } ell > 2;$
 $P \llbracket (ell!0), (ell!1) \rrbracket;$
 $P \llbracket (ell!1)\#(\text{drop } 2 \text{ } ell) \rrbracket \implies P \text{ } ell$
shows $P \text{ } ell$
 $\langle \text{proof} \rangle$

lemma *make-polygonal-path-gives-path:*
fixes $v :: (\text{real}^2) \text{ list}$
shows $\text{path } (\text{make-polygonal-path } v)$
 $\langle \text{proof} \rangle$

corollary *polygonal-path-is-path:*
fixes $g :: R\text{-to-}R^2$
assumes *polygonal-path* g
shows *path* g
 $\langle \text{proof} \rangle$

lemma *polygon-to-polygonal-path:*
fixes $p :: R\text{-to-}R^2$
assumes *polygon* p
obtains ell **where** $p = \text{make-polygonal-path } ell$
 $\langle \text{proof} \rangle$

lemma *polygon-pathstart:*
fixes $g :: R\text{-to-}R^2$
assumes $l \neq []$
assumes $g = \text{make-polygonal-path } l$
shows $\text{pathstart } g = l!0$
 $\langle \text{proof} \rangle$

lemma *polygon-pathfinish:*
fixes $g :: R\text{-to-}R^2$
assumes $l \neq []$
assumes $g = \text{make-polygonal-path } l$
shows $\text{pathfinish } g = l!(\text{length } l - 1)$
 $\langle \text{proof} \rangle$

lemma *make-polygonal-path-image-property:*
assumes $\text{length } vts \geq 2$
assumes *p-is-path:* $x \in \text{path-image } (\text{make-polygonal-path } vts)$
shows $\exists k < \text{length } vts - 1. x \in \text{path-image } (\text{linepath } (vts ! k) (vts ! (k + 1)))$
 $\langle \text{proof} \rangle$

lemma *linepaths-subset-make-polygonal-path-image:*
assumes $\text{length } vts \geq 2$
assumes $k < \text{length } vts - 1$
shows $\text{path-image } (\text{linepath } (vts ! k) (vts ! (k + 1))) \subseteq \text{path-image } (\text{make-polygonal-path } vts)$

<proof>

lemma *vertices-on-path-image*: **shows** $\text{set } vts \subseteq \text{path-image } (\text{make-polygonal-path } vts)$

<proof>

lemma *path-image-cons-union*:

assumes $p = \text{make-polygonal-path } vts$

assumes $p' = \text{make-polygonal-path } vts'$

assumes $vts' \neq []$

assumes $vts = a \# vts' \wedge b = vts!0$

shows $\text{path-image } p = \text{path-image } (\text{linepath } a \ b) \cup \text{path-image } p'$

<proof>

lemma *polygonal-path-image-linepath-union*:

assumes $p = \text{make-polygonal-path } vts$

assumes $n = \text{length } vts$

assumes $n \geq 2$

shows $\text{path-image } p = (\bigcup \{\text{path-image } (\text{linepath } (vts!i) \ (vts!(i+1))) \mid i. i \leq n - 2\})$

<proof>

6 Loop Free Properties

lemma *constant-linepath-is-not-loop-free*:

shows $\neg(\text{loop-free } ((\text{linepath } a \ a)::\text{real} \Rightarrow \text{real}^2))$

<proof>

lemma *doubling-back-is-not-loop-free*:

assumes $a \neq b$

shows $\neg(\text{loop-free } ((\text{make-polygonal-path } [a, b, a])::\text{real} \Rightarrow \text{real}^2))$

<proof>

lemma *not-loop-free-first-component*:

assumes $\neg(\text{loop-free } p1)$

shows $\neg(\text{loop-free } (p1+++p2))$

<proof>

lemma *not-loop-free-second-component*:

assumes $\text{pathfinish-pathstart: pathfinish } p1 = \text{pathstart } p2$

assumes $\neg(\text{loop-free } p2)$

shows $\neg(\text{loop-free } (p1+++p2))$

<proof>

lemma *loop-free-subpath*:

assumes $\text{path } p$

assumes $u\text{-and-}v: u \in \{0..1\} \ v \in \{0..1\} \ u < v$

assumes $\neg(\text{loop-free } (\text{subpath } u \ v \ p))$

shows $\neg(\text{loop-free } p)$

<proof>

lemma *loop-free-associative:*

assumes *path p*

assumes *path q*

assumes *path r*

assumes *pathfinish p = pathstart q*

assumes *pathfinish q = pathstart r*

shows $\neg (\text{loop-free } ((p +++ q) +++ r)) \longleftrightarrow \neg (\text{loop-free } (p +++ (q +++ r)))$

<proof>

lemma *polygon-at-least-3-vertices:*

assumes *polygon p and*

p = make-polygonal-path vts

shows $\text{card } (\text{set } vts) \geq 3$

<proof>

lemma *polygon-vertices-length-at-least-4:*

assumes *polygon p and*

p = make-polygonal-path vts

shows $\text{length } vts \geq 4$

<proof>

lemma *linepath-loop-free:*

assumes $a \neq b$

shows *loop-free (linepath a b)*

<proof>

7 Explicit Linepath Characterization of Polygonal Paths

lemma *triangle-linepath-images:*

fixes $x :: \text{real}$

assumes $vts = [a, b, c]$

assumes $p = \text{make-polygonal-path } vts$

shows $x \in \{0..1/2\} \implies p \ x = ((\text{linepath } a \ b)) \ (2*x)$

$x \in \{1/2..1\} \implies p \ x = ((\text{linepath } b \ c)) \ (2*x - 1)$

<proof>

lemma *polygon-linepath-images1:*

fixes $n :: \text{nat}$

assumes $n \geq 3$

assumes $\text{length } ell = n$

assumes $x \in \{0..1/2\}$

shows $\text{make-polygonal-path } ell \ x = ((\text{linepath } (ell \ ! \ 0) \ (ell \ ! \ 1))) \ (2*x)$

<proof>

lemma *sum-insert* [*simp*]:

assumes $x \notin F$ **and** *finite* F

shows $(\sum y \in \text{insert } x F. P y) = (\sum y \in F. P y) + P x$

<proof>

lemma *sum-of-index-diff* [*simp*]:

fixes $f :: \text{nat} \Rightarrow 'a :: \text{comm-monoid-add}$

shows $(\sum i \in \{a..<a+b\}. f(i-a)) = (\sum i \in \{..<b\}. f(i))$

<proof>

lemma *sum-of-index-diff2* [*simp*]:

fixes $f :: \text{nat} \Rightarrow 'a :: \text{comm-monoid-add}$

shows $(\sum i \in \{a+c..b+c\}. f(i)) = (\sum i \in \{a..b\}. f(i+c))$

<proof>

lemma *sum-split* [*simp*]:

fixes $f :: \text{nat} \Rightarrow 'a :: \text{comm-monoid-add}$

assumes $c \in \{a..b\}$

shows $(\sum i \in \{a..b\}. f i) = (\sum i \in \{a..c\}. f i) + (\sum i \in \{c+1..b\}. f i)$

<proof>

lemma *summation-helper*:

fixes $x :: \text{real}$

fixes $k :: \text{nat}$

assumes $1 \leq k$

shows $(2 :: \text{real}) * (\sum i = 1..k. 1 / 2^i) - 1 = (\sum i = 1..(k-1). (1 / (2^i)))$

<proof>

lemma *polygon-linepath-images2*:

fixes $n k :: \text{nat}$

fixes $ell :: (\text{real}^2) \text{ list}$

fixes $f :: \text{nat} \Rightarrow \text{real} \Rightarrow \text{real}$

assumes $n \geq 3$

assumes $0 \leq k \wedge k \leq n - 3$

assumes $\text{length } ell = n$

assumes $p : p = \text{make-polygonal-path } ell$

assumes $f = (\lambda k x. (x - (\sum i \in \{1..k\}. 1/(2^i))) * (2^{k+1}))$

assumes $x \in \{(\sum i \in \{1..k\}. 1/(2^i))..(\sum i \in \{1..(k+1)\}. 1/(2^i))\}$

shows $p x = ((\text{linepath } (ell ! k) (ell ! (k+1)) (f k x)))$

<proof>

lemma *polygon-linepath-images3*:

fixes $n k :: \text{nat}$

fixes $ell :: (\text{real}^2) \text{ list}$

assumes $n \geq 3$

assumes $\text{length } ell = n$

assumes $p = \text{make-polygonal-path } ell$

assumes $x \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..1\}$

assumes $f = (\lambda x. (x - (\sum i \in \{1..n-2\}. 1/(2^i))) * (2^{n-2}))$
shows $p\ x = (\text{linepath } (ell\ !\ (n-2))\ (ell\ !\ (n-1)))\ (f\ x)$
 <proof>

8 A Triangle is a Polygon

lemma *not-collinear-linepaths-intersect-helper*:

assumes *not-collinear*: $\neg \text{collinear } \{a,b,c\}$
assumes $0 \leq k1$
assumes $k1 \leq 1$
assumes $0 \leq k2$
assumes $k2 \leq 1$
assumes *eo*: $k2 = 0 \implies k1 \neq 1$
shows $\neg ((\text{linepath } a\ b)\ k1 = (\text{linepath } b\ c)\ k2)$
 <proof>

lemma *not-collinear-linepaths-intersect-helper-2*:

assumes *not-collinear*: $\neg \text{collinear } \{a,b,c\}$
assumes $0 \leq k1$
assumes $k1 \leq 1$
assumes $0 \leq k2$
assumes $k2 \leq 1$
assumes *eo*: $k1 = 0 \implies k2 \neq 1$
shows $\neg ((\text{linepath } a\ b)\ k1 = (\text{linepath } c\ a)\ k2)$
 <proof>

lemma *not-collinear-loopfree-path*: $\bigwedge a\ b\ c::\text{real}^2. \neg \text{collinear } \{a,b,c\} \implies \text{loop-free } ((\text{linepath } a\ b) \text{ +++ } (\text{linepath } b\ c))$

<proof>

lemma *triangle-is-polygon*: $\bigwedge a\ b\ c. \neg \text{collinear } \{a,b,c\} \implies \text{polygon } (\text{make-triangle } a\ b\ c)$

<proof>

lemma *have-wraparound-vertex*:

assumes *polygon* p
assumes $p = \text{make-polygonal-path } vts$
shows $vts = (\text{take } (\text{length } vts - 1)\ vts)@[vts\ !\ 0]$
 <proof>

lemma *polygon-at-least-3-vertices-wraparound*:

assumes *polygon* p
assumes $p = \text{make-polygonal-path } vts$
shows $\text{card } (\text{set } (\text{take } (\text{length } vts - 1)\ vts)) \geq 3$
 <proof>

9 Polygon Vertex Rotation

definition *rotate-polygon-vertices*:: 'a list \Rightarrow nat \Rightarrow 'a list
where *rotate-polygon-vertices* ell i =
 (let ell1 = rotate i (butlast ell) in ell1 @ [ell1 ! 0])

lemma *rotate-polygon-vertices-same-set*:
assumes polygon (make-polygonal-path vts)
shows set (rotate-polygon-vertices vts i) = set vts
 <proof>

lemma *arb-rotation-as-single-rotation*:
fixes i:: nat
shows rotate-polygon-vertices vts (Suc i) = rotate-polygon-vertices (rotate-polygon-vertices vts i) 1
 <proof>

lemma *rotation-sum*:
fixes i j :: nat
shows rotate-polygon-vertices vts (i + j) = rotate-polygon-vertices (rotate-polygon-vertices vts i) j
 <proof>

lemma *rotated-polygon-vertices-helper*:
fixes p :: R-to-R2
assumes poly-p: polygon p
assumes p-is-path: p = make-polygonal-path vts
assumes p'-is: p' = make-polygonal-path (rotate-polygon-vertices vts 1)
shows (vts ! 0) = (rotate-polygon-vertices vts 1) ! (length (rotate-polygon-vertices vts 1) - 2)
 (rotate-polygon-vertices vts 1) ! (length (rotate-polygon-vertices vts 1) - 1)
 = (vts ! 1)
 <proof>

lemma *rotate-polygon-vertices-same-length*:
fixes vts :: (real²) list
assumes length vts \geq 1
shows length vts = length (rotate-polygon-vertices vts i)
 <proof>

lemma *rotated-polygon-vertices-helper2*:
assumes len-gteq: length vts \geq 2
assumes i < length vts - 1
assumes hd vts = last vts
shows (rotate-polygon-vertices vts 1) ! i = vts ! (i+1)
 <proof>

lemma *polygon-rotation-t-translation1*:
assumes polygon-of p vts

assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
 (is $p' = \text{make-polygonal-path } ?vts'$)
assumes $x' \in \{(\sum i \in \{1..k\}. 1/(2^{\widehat{i}}))..(\sum i \in \{1..k+1\}. 1/(2^{\widehat{i}}))\}$
assumes $n = \text{length } vts$
assumes $0 \leq k \wedge k \leq n - 4$
assumes $l = x' - (\sum i \in \{1..k\}. 1/(2^{\widehat{i}}))$
assumes $x = l/2 + (\sum i \in \{1..(k+1)\}. 1/(2^{\widehat{i}}))$
shows $x \in \{(\sum i \in \{1..k+1\}. 1/(2^{\widehat{i}}))..(\sum i \in \{1..k+2\}. 1/(2^{\widehat{i}}))\}$
 $p' x' = p x$
 <proof>

lemma *polygon-rotation-t-translation1-strict:*

assumes *polygon-of p vts*
assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
 (is $p' = \text{make-polygonal-path } ?vts'$)
assumes $x' \in \{(\sum i \in \{1..k\}. 1/(2^{\widehat{i}}))..<(\sum i \in \{1..k+1\}. 1/(2^{\widehat{i}}))\}$
assumes $n = \text{length } vts$
assumes $0 \leq k \wedge k \leq n - 4$
assumes $l = x' - (\sum i \in \{1..k\}. 1/(2^{\widehat{i}}))$
assumes $x = l/2 + (\sum i \in \{1..(k+1)\}. 1/(2^{\widehat{i}}))$
shows $x \in \{(\sum i \in \{1..k+1\}. 1/(2^{\widehat{i}}))..<(\sum i \in \{1..k+2\}. 1/(2^{\widehat{i}}))\}$
 $p' x' = p x$
 <proof>

lemma *polygon-rotation-t-translation2:*

assumes *polygon-of p vts*
assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
 (is $p' = \text{make-polygonal-path } ?vts'$)
assumes $n = \text{length } vts$
assumes $x' \in \{(\sum i \in \{1..(n-3)\}. 1/(2^{\widehat{i}}))..(\sum i \in \{1..(n-2)\}. 1/(2^{\widehat{i}}))\}$
assumes $x = x' + 1/(2^{\widehat{(n-2)}})$
shows $x \in \{(\sum i \in \{1..n-2\}. 1/(2^{\widehat{i}}))..1\}$
 $p' x' = p x$
 <proof>

lemma *polygon-rotation-t-translation2-strict:*

assumes *polygon-of p vts*
assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
 (is $p' = \text{make-polygonal-path } ?vts'$)
assumes $n = \text{length } vts$
assumes $x' \in \{(\sum i \in \{1..(n-3)\}. 1/(2^{\widehat{i}}))..<(\sum i \in \{1..(n-2)\}. 1/(2^{\widehat{i}}))\}$
assumes $x = x' + 1/(2^{\widehat{(n-2)}})$
shows $x \in \{(\sum i \in \{1..n-2\}. 1/(2^{\widehat{i}}))..<1\}$
 $p' x' = p x$
 <proof>

lemma *polygon-rotation-t-translation3:*

assumes *polygon-of p vts*

assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
 (is $p' = \text{make-polygonal-path } ?vts'$)
assumes $x' \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..1\}$
assumes $n = \text{length } vts$
assumes $l = x' - (\sum i \in \{1..n-2\}. 1/(2^i))$
assumes $x = l * (2^{n-3})$
shows $x \in \{0..1/2\}$
 $p' \ x' = p \ x$
 <proof>

lemma *polygon-rotation-t-translation3-strict*:
assumes *polygon-of* $p \ vts$
assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
 (is $p' = \text{make-polygonal-path } ?vts'$)
assumes $x' \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..<1\}$
assumes $n = \text{length } vts$
assumes $l = x' - (\sum i \in \{1..n-2\}. 1/(2^i))$
assumes $x = l * (2^{n-3})$
shows $x \in \{0..<1/2\}$
 $p' \ x' = p \ x$
 <proof>

lemma *f-gteq-0-sum-gt*: $\bigwedge f::nat \Rightarrow real. (\bigwedge i::nat. (f \ i) > 0) \Longrightarrow a > b \Longrightarrow (\sum i = 1..a. (f \ i)) > (\sum i = 1..b. (f \ i))$ for $a \ b :: nat$
 <proof>

lemma *rotation-intervals-disjoint*:
assumes $k1 \neq k2$
shows $\{\sum i = 1..k1. 1 / (2^i::real)..<\sum i = 1..k1+1. 1 / 2^i\} \cap \{\sum i = 1..k2. 1 / (2^i::real)..<\sum i = 1..k2+1. 1 / 2^i\} = \{\}$
 <proof>

lemma *bounding-interval-helper1*:
shows $(\sum i = 1..k. 1 / (2^i::real)) = (2^k - 1)/(2^k)$
 <proof>

lemma *bounding-interval-helper2*:
fixes $x :: real$
assumes $x \in \{0..<1\}$
shows $\exists k. x < (\sum i = 1..k. 1 / (2^i::real))$
 <proof>

lemma *bounding-interval-for-reals-btw01*:
fixes $x::real$
assumes $x \in \{0..<1\}$
shows $\exists k. x \in \{(\sum i \in \{1..k\}. 1/(2^i::real))..<(\sum i \in \{1..(k+1)\}. 1/(2^i))\}$
 <proof>

lemma *all-rotation-intervals-between-0and1*:

shows $\{(\sum i \in \{1..k\}. 1/(2^{\widehat{i}}::real))..(\sum i \in \{1..(k+1)\}. 1/(2^{\widehat{i}}))\} \subseteq \{0..<1\}$
<proof>

lemma *all-rotation-intervals-between-0and1-strict:*

shows $\{(\sum i \in \{1..k\}. 1/(2^{\widehat{i}}::real))..<(\sum i \in \{1..(k+1)\}. 1/(2^{\widehat{i}}))\} \subseteq \{0..<1\}$
<proof>

lemma *one-polygon-rotation-is-loop-free:*

assumes *polygon-of p vts*

assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
(**is** $p' = \text{make-polygonal-path } ?vts'$)

shows *loop-free p'*

<proof>

lemma *one-rotation-is-polygon:*

fixes $p :: R\text{-to-}R^2$

fixes $i :: nat$

assumes *poly-p: polygon p and*

p-is-path: p = make-polygonal-path vts and

p'-is: p' = make-polygonal-path (rotate-polygon-vertices vts 1)

(**is** $p' = \text{make-polygonal-path } ?vts'$)

shows *polygon p'*

<proof>

lemma *rotation-is-polygon:*

fixes $p :: R\text{-to-}R^2$

fixes $i :: nat$

assumes *polygon p and*

p = make-polygonal-path vts

shows *polygon (make-polygonal-path (rotate-polygon-vertices vts i))*

<proof>

lemma *polygon-rotate-mod:*

fixes $vts :: (real^2) \text{ list}$

assumes $n = \text{length } vts$

assumes $n \geq 2$

assumes $\text{hd } vts = \text{last } vts$

shows *rotate-polygon-vertices vts (n - 1) = vts*

<proof>

lemma *polygon-rotate-mod-arb:*

fixes $vts :: (real^2) \text{ list}$

assumes $n = \text{length } vts$

assumes $n \geq 2$

assumes $\text{hd } vts = \text{last } vts$

shows *rotate-polygon-vertices vts ((n - 1) * i) = vts*

<proof>

lemma *unrotation-is-polygon:*

fixes $p :: R\text{-to-}R^2$
fixes $i :: \text{nat}$
assumes $\text{polygon } (\text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ i))$
 $(\text{is polygon } (\text{make-polygonal-path } ?vts'))$
 $p = \text{make-polygonal-path } vts$
 $\text{hd } vts = \text{last } vts$
shows $\text{polygon } p$
 $\langle \text{proof} \rangle$

lemma $\text{rotated-polygon-vertices}$:
assumes $vts' = \text{rotate-polygon-vertices } vts \ j$
assumes $\text{hd } vts = \text{last } vts$
assumes $\text{length } vts \geq 2$
assumes $j \leq i \wedge i < \text{length } vts$
shows $vts \ ! \ i = vts' \ ! \ (i - j)$
 $\langle \text{proof} \rangle$

lemma $\text{polygon-path-image}$:
assumes $\text{poly-p: polygon } p$
assumes $\text{p-is-path: } p = \text{make-polygonal-path } vts$
shows $\text{path-image } p = p' \ \{0 \ .. < \ 1\}$
 $\langle \text{proof} \rangle$

lemma $\text{polygon-vts-one-rotation}$:
fixes $p :: R\text{-to-}R^2$
assumes $\text{poly-p: polygon } p$ **and**
 $\text{p-is-path: } p = \text{make-polygonal-path } vts$ **and**
 $\text{p'-is: } p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
shows $\text{path-image } p = \text{path-image } p'$
 $\langle \text{proof} \rangle$

lemma $\text{polygon-vts-arb-rotation}$:
fixes $p :: R\text{-to-}R^2$
assumes $\text{polygon } p$ **and**
 $p = \text{make-polygonal-path } vts$
shows $\text{path-image } p = \text{path-image } (\text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ i))$
 $\langle \text{proof} \rangle$

10 Translating a Polygon

lemma $\text{linepath-translation}$:
 $\text{linepath } ((\lambda x. x + u) \ a) \ ((\lambda x. x + u) \ b) = (\lambda x. x + u) \circ (\text{linepath } a \ b)$
 $\langle \text{proof} \rangle$

lemma $\text{make-polygonal-path-translate}$:
assumes $\text{length } vts \geq 2$
shows $\text{make-polygonal-path } (\text{map } (\lambda x. x + u) \ vts) = (\lambda x. x + u) \circ (\text{make-polygonal-path } vts)$

<proof>

lemma *translation-is-polygon*:

assumes *polygon-of* p pts

shows *polygon-of* $((\lambda x. x + u) \circ p)$ $(\text{map } (\lambda x. x + u) pts)$ **(is** *polygon-of* $?p'$ $?pts'$)

<proof>

11 Misc. properties

lemma *tail-of-loop-free-polygonal-path-is-loop-free*:

assumes *loop-free* $(\text{make-polygonal-path } (x\#\text{tail}))$ **(is** *loop-free* $?p$) **and**
 $\text{length tail} \geq 2$

shows *loop-free* $(\text{make-polygonal-path tail})$ **(is** *loop-free* $?p'$)

<proof>

lemma *tail-of-simple-polygonal-path-is-simple*:

assumes *simple-path* $(\text{make-polygonal-path } (x\#\text{tail}))$ **(is** *simple-path* $?p$) **and**
 $\text{length tail} \geq 2$

shows *simple-path* $(\text{make-polygonal-path tail})$ **(is** *simple-path* $?p'$)

<proof>

lemma *interior-vertex-in-path-image-interior*:

fixes $pts :: (\text{real}^2)$ *list*

assumes $x \in \text{set } (\text{butlast } (\text{drop } 1 pts))$

shows $\exists t. t \in \{0 < .. < 1\} \wedge (\text{make-polygonal-path } pts) t = x$

<proof>

lemma *loop-free-polygonal-path-pts-distinct*:

assumes *loop-free* $(\text{make-polygonal-path } pts)$

shows *distinct* $(\text{butlast } pts)$

<proof>

lemma *loop-free-polygonal-path-pts-drop1-distinct*:

assumes *loop-free* $(\text{make-polygonal-path } pts)$

shows *distinct* $(\text{drop } 1 pts)$

<proof>

lemma *simple-polygonal-path-pts-distinct*:

assumes *simple-path* $(\text{make-polygonal-path } pts)$

shows *distinct* $(\text{butlast } pts)$

<proof>

lemma *edge-subset-path-image*:

assumes $p = \text{make-polygonal-path } pts$ **and**

$(i::\text{int}) \in \{0..<((\text{length } pts) - 1)\}$ **and**

$x = pts!i$ **and**

$y = vts!(i+1)$
shows $path\text{-}image\ (linepath\ x\ y) \subseteq path\text{-}image\ p$ (**is** $?xy\text{-}img \subseteq ?p\text{-}img$)
 $\langle proof \rangle$

12 Properties of Sublists of Polygonal Path Vertex Lists

lemma *make-polygonal-path-image-append-var:*

assumes $length\ vts1 \geq 2$
shows $path\text{-}image\ (make\text{-}polygonal\text{-}path\ (vts1\ @\ [v])) = path\text{-}image\ (make\text{-}polygonal\text{-}path\ vts1\ +++\ (linepath\ (vts1\ !\ (length\ vts1\ -\ 1))\ v))$
 $\langle proof \rangle$

lemma *make-polygonal-path-image-append-helper:*

assumes $length\ vts1 \geq 1 \wedge length\ vts2 \geq 1$
shows $path\text{-}image\ (make\text{-}polygonal\text{-}path\ (vts1\ @\ [v]\ @\ [v]\ @\ vts2)) = path\text{-}image\ (make\text{-}polygonal\text{-}path\ (vts1\ @\ [v]\ @\ vts2))$
 $\langle proof \rangle$

lemma *make-polygonal-path-image-append:*

assumes $length\ vts1 \geq 2 \wedge length\ vts2 \geq 2$
shows $path\text{-}image\ (make\text{-}polygonal\text{-}path\ (vts1\ @\ vts2)) = path\text{-}image\ (make\text{-}polygonal\text{-}path\ vts1\ +++\ (linepath\ (vts1\ !\ (length\ vts1\ -\ 1))\ (vts2\ !\ 0))\ +++\ make\text{-}polygonal\text{-}path\ vts2)$
 $\langle proof \rangle$

lemma *make-polygonal-path-image-append-alt:*

assumes $p = make\text{-}polygonal\text{-}path\ vts$
assumes $p1 = make\text{-}polygonal\text{-}path\ vts1$
assumes $p2 = make\text{-}polygonal\text{-}path\ vts2$
assumes $last\ vts1 = hd\ vts2$
assumes $length\ vts1 \geq 2 \wedge length\ vts2 \geq 2$
assumes $vts = vts1\ @\ (tl\ vts2)$
shows $path\text{-}image\ p = path\text{-}image\ (p1\ +++\ p2)$
 $\langle proof \rangle$

lemma *cont-incr-interval-image:*

fixes $f :: real \Rightarrow real$
assumes $a \leq b$
assumes *continuous-on* $\{a..b\}\ f$
assumes $\forall x \in \{a..b\}. \forall y \in \{a..b\}. x \leq y \longrightarrow f\ x \leq f\ y$
shows $f'\{a..b\} = \{f\ a..f\ b\}$
 $\langle proof \rangle$

lemma *two-x-minus-one-image:*

assumes $f = (\lambda x :: real. 2*x - 1)$
assumes $a \leq b$
shows $f'\{a..b\} = \{f\ a..f\ b\}$

<proof>

lemma *vts-split-path-image:*

assumes $p = \text{make-polygonal-path } vts$
assumes $p1 = \text{make-polygonal-path } vts1$
assumes $p2 = \text{make-polygonal-path } vts2$
assumes $vts1 = \text{take } i \text{ } vts$
assumes $vts2 = \text{drop } (i-1) \text{ } vts$
assumes $n = \text{length } vts$
assumes $1 \leq i \wedge i < n$
assumes $x = (2^{i-1} - 1) / (2^{i-1})$
shows $\text{path-image } p1 = p\{0..x\} \wedge \text{path-image } p2 = p\{x..1\}$
<proof>

lemma *drop-i-is-loop-free:*

fixes $vts :: (\text{real}^2) \text{ list}$
assumes $m = \text{length } vts$
assumes $i \leq m - 2$
assumes $vts' = \text{drop } i \text{ } vts$
assumes $p = \text{make-polygonal-path } vts$
assumes $p' = \text{make-polygonal-path } vts'$
assumes *loop-free* p
shows *loop-free* p'
<proof>

lemma *joinpaths-tl-transform:*

assumes $f = (\lambda x :: \text{real}. 2*x - 1)$
assumes $\text{pathfinish } g1 = \text{pathstart } g2$
assumes $p = g1 \text{ +++ } g2$
assumes $x \geq 1/2$
shows $p \text{ } x = g2 \text{ } (f \text{ } x)$
<proof>

lemma *joinpaths-tl-image-transform:*

assumes $f = (\lambda x :: \text{real}. 2*x - 1)$
assumes $\text{pathfinish } g1 = \text{pathstart } g2$
assumes $p = g1 \text{ +++ } g2$
assumes $1/2 \leq a \wedge a \leq b$
shows $p\{a..b\} = g2\{f \text{ } a..f \text{ } b\}$
<proof>

lemma *vts-sublist-path-image:*

assumes $p = \text{make-polygonal-path } vts$
assumes $p' = \text{make-polygonal-path } vts'$
assumes $vts' = \text{take } j \text{ } (\text{drop } i \text{ } vts)$
assumes $m = \text{length } vts$
assumes $n = \text{length } vts'$
assumes $k = i + j$
assumes $k \leq m - 1 \wedge 2 \leq j$

assumes $x1 = (2^i - 1)/(2^i)$
assumes $x2 = (2^{k-1} - 1)/(2^{k-1})$
shows $\text{path-image } p' = p'\{x1..x2\}$
 <proof>

lemma *one-append-simple-path:*

fixes $pts :: (\mathbb{R}^2) \text{ list}$
assumes $pts = pts' @ [z]$
assumes $n = \text{length } pts$
assumes $n \geq 3$
assumes $p = \text{make-polygonal-path } pts$
assumes $p' = \text{make-polygonal-path } pts'$
assumes $\text{simple-path } p$
shows $\text{simple-path } p'$
 <proof>

lemma *take-i-is-loop-free:*

fixes $pts :: (\mathbb{R}^2) \text{ list}$
assumes $n = \text{length } pts$
assumes $2 \leq i \wedge i \leq n$
assumes $pts' = \text{take } i \text{ } pts$
assumes $p = \text{make-polygonal-path } pts$
assumes $p' = \text{make-polygonal-path } pts'$
assumes $\text{loop-free } p$
shows $\text{loop-free } p'$
 <proof>

lemma *sublist-is-loop-free:*

fixes $pts :: (\mathbb{R}^2) \text{ list}$
assumes $p = \text{make-polygonal-path } pts$
assumes $p' = \text{make-polygonal-path } pts'$
assumes $\text{loop-free } p$
assumes $m = \text{length } pts$
assumes $n = \text{length } pts'$
assumes $\text{sublist } pts' \text{ } pts$
assumes $n \geq 2 \wedge m \geq 2$
shows $\text{loop-free } p'$
 <proof>

lemma *diff-points-path-image-set-property:*

fixes $a \ b :: \mathbb{R}^2$
assumes $a \neq b$
shows $\text{path-image } (\text{linepath } a \ b) \neq \{a, b\}$
 <proof>

lemma *polygonal-path-vertex-t:*

assumes $p = \text{make-polygonal-path } pts$
assumes $n = \text{length } pts$
assumes $n \geq 1$

assumes $0 \leq i \wedge i < n - 1$
assumes $x = (2^{\widehat{i}} - 1) / (2^{\widehat{i}})$
shows $vts!i = p \ x$
 <proof>

lemma *loop-free-split-int:*

assumes $p = \text{make-polygonal-path } vts \wedge \text{loop-free } p$
assumes $vts1 = \text{take } i \ vts$
assumes $vts2 = \text{drop } (i-1) \ vts$
assumes $c1 = \text{make-polygonal-path } vts1$
assumes $c2 = \text{make-polygonal-path } vts2$
assumes $n = \text{length } vts$
assumes $1 \leq i \wedge i < n$
shows $(\text{path-image } c1) \cap (\text{path-image } c2) \subseteq \{\text{pathstart } c1, \text{pathstart } c2\}$
 (is $?C1 \cap ?C2 \subseteq \{\text{pathstart } c1, \text{pathstart } c2\}$)
 <proof>

lemma *loop-free-arc-split-int:*

assumes $p = \text{make-polygonal-path } vts \wedge \text{loop-free } p \wedge \text{arc } p$
assumes $vts1 = \text{take } i \ vts$
assumes $vts2 = \text{drop } (i-1) \ vts$
assumes $c1 = \text{make-polygonal-path } vts1$
assumes $c2 = \text{make-polygonal-path } vts2$
assumes $n = \text{length } vts$
assumes $1 \leq i \wedge i < n$
shows $(\text{path-image } c1) \cap (\text{path-image } c2) \subseteq \{\text{pathstart } c2\}$
 (is $?C1 \cap ?C2 \subseteq \{\text{pathstart } c2\}$)
 <proof>

lemma *loop-free-append:*

assumes $p = \text{make-polygonal-path } vts$
assumes $p1 = \text{make-polygonal-path } vts1$
assumes $p2 = \text{make-polygonal-path } vts2$
assumes $vts = vts1 \ @ \ (\text{tl } vts2)$
assumes $\text{loop-free } p1 \wedge \text{loop-free } p2$
assumes $\text{path-image } p1 \cap \text{path-image } p2 \subseteq \{\text{pathstart } p1, \text{pathstart } p2\}$
assumes $\text{last } vts2 \neq \text{hd } vts1 \longrightarrow \text{path-image } p1 \cap \text{path-image } p2 \subseteq \{\text{pathstart } p2\}$
assumes $\text{last } vts1 = \text{hd } vts2$
assumes $\text{arc } p1 \wedge \text{arc } p2$
shows $\text{loop-free } p$
 <proof>

lemma *sublist-path-image-subset:*

assumes $\text{sublist } vts1 \ vts2$
assumes $\text{length } vts1 \geq 1$
shows $\text{path-image } (\text{make-polygonal-path } vts1) \subseteq \text{path-image } (\text{make-polygonal-path } vts2)$
 <proof>

lemma *integral-on-edge-subset-integral-on-path*:
assumes $p = \text{make-polygonal-path } vts$ **and**
 $(i::int) \in \{0..<((\text{length } vts) - 1)\}$ **and**
 $x = vts!i$ **and**
 $y = vts!(i+1)$
shows $\{v. \text{integral-vec } v \wedge v \in \text{path-image } (\text{linepath } x \ y)\}$
 $\subseteq \{v. \text{integral-vec } v \wedge v \in \text{path-image } p\}$
 $\langle \text{proof} \rangle$

lemma *sublist-pair-integral-subset-integral-on-path*:
assumes $p = \text{make-polygonal-path } vts$ **and**
 $\text{sublist } [x, y] \ vts$
shows $\{v. \text{integral-vec } v \wedge v \in \text{path-image } (\text{linepath } x \ y)\}$
 $\subseteq \{v. \text{integral-vec } v \wedge v \in \text{path-image } p\}$
 $\langle \text{proof} \rangle$

lemma *sublist-integral-subset-integral-on-path*:
assumes $\text{length } ell \geq 2$
assumes $p = \text{make-polygonal-path } vts$ **and**
 $\text{sublist } ell \ vts$
shows $\{v. \text{integral-vec } v \wedge v \in \text{path-image } (\text{make-polygonal-path } ell)\}$
 $\subseteq \{v. \text{integral-vec } v \wedge v \in \text{path-image } p\}$
 $\langle \text{proof} \rangle$

13 Reversing Polygonal Path Vertex List

lemma *rev-vts-path-image*:
shows $\text{path-image } (\text{make-polygonal-path } (\text{rev } vts)) = \text{path-image } (\text{make-polygonal-path } vts)$
 $\langle \text{proof} \rangle$

lemma *rev-vts-is-loop-free*:
assumes $p = \text{make-polygonal-path } vts$
assumes $\text{loop-free } p$
shows $\text{loop-free } (\text{make-polygonal-path } (\text{rev } vts))$
 $\langle \text{proof} \rangle$

lemma *rev-vts-is-polygon*:
assumes $\text{polygon-of } p \ vts$
shows $\text{polygon } (\text{make-polygonal-path } (\text{rev } vts))$
 $\langle \text{proof} \rangle$

end
theory *Linepath-Collinearity*
imports *Polygon-Lemmas*

begin

14 Collinearity Properties

lemma *points-on-linepath-collinear*:

assumes *exists-c*: $(\exists c. a - b = c *_R u)$

assumes *x-in-linepath*: $x \in \text{path-image } (\text{linepath } a \ b)$

shows $(\exists c. x - a = c *_R u) (\exists c. b - x = c *_R u)$

<proof>

lemma *three-points-collinear-property*:

fixes *a b*:: real^2

assumes *exists-c1*: $(\exists c. a - x1 = c *_R u)$

assumes *exists-c2*: $(\exists c. a - x2 = c *_R u)$

shows $\exists c. x1 - x2 = c *_R u$

<proof>

lemma *in-path-image-imp-collinear*:

fixes *a b*:: real^2

assumes *k* $\in \text{path-image } (\text{linepath } a \ b)$

shows *collinear* $\{a, b, k\}$

<proof>

lemma *two-linepath-collinearity-property*:

fixes *a b c d*:: real^2

assumes $y \neq z \wedge \{y, z\} \subseteq (\text{path-image } (\text{linepath } a \ b)) \cap (\text{path-image } (\text{linepath } c \ d))$

shows *collinear* $\{a, b, c, d\}$

<proof>

lemma *polygon-vts-not-collinear*:

assumes *polygon-of p vts*

shows $\neg \text{collinear } (\text{set } vts)$

<proof>

lemma *not-collinear-with-subset*:

assumes *collinear A*

assumes $\neg \text{collinear } (A \cup \{x\})$

assumes $\text{card } A > 2$

assumes $a \in A$

shows $\neg \text{collinear } ((A - \{a\}) \cup \{x\})$

<proof>

lemma *vec-diff-scale-collinear*:

fixes *a b c* :: real^2

assumes $b - a = m *_R (c - a)$

shows *collinear* $\{a, b, c\}$

<proof>

15 Linepath Properties

lemma *good-linepath-comm*: $\text{good-linepath } a \ b \ vts \implies \text{good-linepath } b \ a \ vts$
 ⟨proof⟩

lemma *finite-set-linepaths*:
 assumes *polygon*: $\text{polygon } p$
 assumes *polygonal-path*: $p = \text{make-polygonal-path } vts$
 shows *finite* $\{(a, b). (a, b) \in \text{set } vts \times \text{set } vts\}$
 ⟨proof⟩

lemma *linepaths-intersect-once-or-collinear*:
 fixes $a \ b \ c \ d :: \text{real}^2$
 assumes *path-image* $(\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d) \neq \{\}$
 shows *collinear* $\{a, b, c, d\} \vee (\exists x. \text{path-image } (\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d) = \{x\})$
 ⟨proof⟩

lemma *linepaths-intersect-once-or-collinear-alt*:
 fixes $a \ b \ c \ d :: \text{real}^2$
 assumes *path-image* $(\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d) \neq \{\}$
 shows *collinear* $\{a, b, c, d\} \vee \text{card } (\text{path-image } (\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d)) = 1$
 ⟨proof⟩

lemma *path-image-linepath-union*:
 fixes $a \ b :: 'a :: \text{euclidean-space}$
 assumes $d \in \text{path-image } (\text{linepath } a \ b)$
 shows $\text{path-image } (\text{linepath } a \ b) = \text{path-image } (\text{linepath } a \ d) \cup \text{path-image } (\text{linepath } d \ b)$
 ⟨proof⟩

lemma *path-image-linepath-split*:
 assumes $i < (\text{length } vts) - 1$
 assumes $x \in \text{path-image } (\text{linepath } (vts!i) \ (vts!(i+1)))$
 assumes *x-notin*: $x \notin \text{set } vts$
 shows $\text{path-image } (\text{make-polygonal-path } vts) = \text{path-image } (\text{make-polygonal-path } ((\text{take } (i+1) \ vts) @ [x] @ (\text{drop } (i+1) \ vts)))$
 ⟨proof⟩

lemma *linepath-split-is-loop-free*:
 assumes $d \in \text{path-image } (\text{linepath } a \ b)$
 assumes $d \notin \{a, b\}$
 shows *loop-free* $(\text{make-polygonal-path } [a, d, b])$ (is loop-free ?p)
 ⟨proof⟩

lemma *loop-free-linepath-split-is-loop-free*:
 assumes $p = \text{make-polygonal-path } vts$
 assumes *loop-free* p

assumes $n = \text{length } vts$
assumes $i < n - 1$
assumes $x \in \text{path-image } (\text{linepath } (vts!i) (vts!(i+1))) \wedge x \notin \text{set } vts$
assumes $vts' = (\text{take } (i+1) vts) @ [x] @ (\text{drop } (i+1) vts)$
assumes $p' = \text{make-polygonal-path } vts'$
shows $\text{loop-free } p' \wedge \text{path-image } p' = \text{path-image } p$
 <proof>

lemma *polygon-linepath-split-is-polygon*:
assumes *polygon-of* p vts
assumes $i < (\text{length } vts) - 1$
assumes $a = vts!i \wedge b = vts!(i+1)$
assumes $x \in \text{path-image } (\text{linepath } a b) \wedge x \notin \text{set } vts$
assumes $vts' = (\text{take } (i+1) vts) @ [x] @ (\text{drop } (i+1) vts)$
shows *polygon* ($\text{make-polygonal-path } vts'$)
 <proof>

16 Measure of linepaths

lemma *linepath-is-negligible-vertical*:
fixes $a b :: \text{real}^2$
assumes $a\$1 = b\1
defines $p \equiv \text{linepath } a b$
shows *negligible* ($\text{path-image } p$)
 <proof>

lemma *linepath-is-negligible-non-vertical*:
fixes $a b :: \text{real}^2$
assumes $a\$1 < b\1
defines $p \equiv \text{linepath } a b$
shows *negligible* ($\text{path-image } p$)
 <proof>

lemma *linepath-is-negligible*:
fixes $a b :: \text{real}^2$
defines $p \equiv \text{linepath } a b$
shows *negligible* ($\text{path-image } p$)
 <proof>

lemma *linepath-has-emeasure-0*:
 $\text{emeasure lebesgue } (\text{path-image } (\text{linepath } (a::(\text{real}^2)) (b::(\text{real}^2)))) = 0$
 <proof>

lemma *linepath-has-measure-0*:
 $\text{measure lebesgue } (\text{path-image } (\text{linepath } (a::(\text{real}^2)) (b::(\text{real}^2)))) = 0$
 <proof>

end

theory *Polygon-Convex-Lemmas*

imports

Polygon-Lemmas

Linepath-Collinearity

begin

17 Misc. Convex Polygon Properties

lemma *polygon-path-image-subset-convex:*

assumes $\text{length } vts > 0$

shows $\text{path-image } (\text{make-polygonal-path } vts) \subseteq \text{convex hull } (\text{set } vts)$ (**is** $\text{path-image } ?p \subseteq ?S$)

<proof>

lemma *convex-contains-simple-closed-path-imp-contains-path-inside:*

assumes $\text{convex } S$

assumes $\text{simple-path } p \wedge \text{closed-path } p$

assumes $\text{path-image } p \subseteq S$

shows $\text{path-inside } p \subseteq S$

<proof>

lemma *convex-polygon-is-convex-hull:*

assumes $\text{polygon } p$

assumes $\text{convex } (\text{path-inside } p \cup \text{path-image } p)$

assumes $p = \text{make-polygonal-path } vts$

shows $\text{convex hull } (\text{set } vts) = \text{path-inside } p \cup \text{path-image } p$ (**is** $?hull = ?poly$)

<proof>

lemma *convex-polygon-inside-is-convex-hull-interior:*

assumes $\text{polygon } p$

assumes $\text{convex } (\text{path-inside } p)$

assumes $p = \text{make-polygonal-path } vts$

shows $\text{interior } (\text{convex hull } (\text{set } vts)) = \text{path-inside } p$

<proof>

lemma *convex-polygon-inside-is-convex-hull-interior2:*

assumes $\text{polygon } p$

assumes $\text{convex } (\text{path-inside } p \cup \text{path-image } p)$

assumes $p = \text{make-polygonal-path } vts$

shows $\text{interior } (\text{convex hull } (\text{set } vts)) = \text{path-inside } p$

<proof>

lemma *polygon-convex-iff:*

assumes $\text{polygon } p$

shows $\text{convex } (\text{path-inside } p) \iff \text{convex } (\text{path-inside } p \cup \text{path-image } p)$

<proof>

lemma *convex-polygon-frontier-is-path-image:*

assumes *polygon-of p vts*
assumes *convex (path-inside p)*
shows $\text{frontier } (\text{convex hull } (\text{set vts})) = \text{path-image } p$
<proof>

lemma *convex-polygon-frontier-is-path-image2:*
assumes *polygon p*
assumes *convex (path-inside p)*
shows $\text{frontier } (\text{path-image } p \cup \text{path-inside } p) = \text{path-image } p$
<proof>

lemma *convex-polygon-frontier-is-path-image3:*
assumes *polygon p*
assumes *convex (path-image p \cup path-inside p)*
shows $\text{frontier } (\text{path-image } p \cup \text{path-inside } p) = \text{path-image } p$
<proof>

lemma *polygon-frontier-is-path-image:*
assumes *polygon p*
shows $\text{frontier } (\text{path-inside } p) = \text{path-image } p$
<proof>

lemma *convex-path-inside-means-convex-polygon:*
assumes *polygon p*
assumes $\text{frontier } (\text{convex hull } (\text{set vts})) = \text{path-image } p$
shows *convex (path-inside p)*
<proof>

lemma *convex-hull-of-polygon-is-convex-hull-of-vts:*
assumes *polygon-of p vts*
shows $\text{convex hull } (\text{path-image } p \cup \text{path-inside } p) = \text{convex hull } (\text{set vts})$
<proof>

lemma *convex-hull-frontier-polygon:*
assumes *polygon-of p vts*
assumes $\neg \text{set vts} \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$
shows $\neg \text{convex } (\text{path-inside } p)$
<proof>

lemma *frontier-int-subset:*
assumes $A \subseteq B$
shows $(\text{frontier } B) \cap A \subseteq \text{frontier } A$
<proof>

lemma *in-frontier-in-subset:*
assumes $A \subseteq B$
assumes $x \in \text{frontier } B$
assumes $x \in A$
shows $x \in \text{frontier } A$

<proof>

lemma *in-frontier-in-subset-convex-hull:*

assumes $A \subseteq B$

assumes $x \in \text{frontier} (\text{convex hull } B)$

assumes $x \in \text{convex hull } A$

shows $x \in \text{frontier} (\text{convex hull } A)$

<proof>

lemma *convex-hull-two-extreme-points:*

fixes $S :: 'a::\text{euclidean-space set}$

assumes *finite* S

assumes $\text{convex hull } S \neq \{\}$

assumes $\forall x. \text{convex hull } S \neq \{x\}$

shows $\text{card } \{x. x \text{ extreme-point-of } (\text{convex hull } S)\} \geq 2$ (**is** $\text{card } ?ep \geq 2$)

<proof>

lemma *convex-hull-two-vts-on-frontier:*

fixes $S :: 'a::\text{euclidean-space set}$

assumes $\text{card } S \geq 2$

shows $\text{card } (S \cap \text{frontier} (\text{convex hull } S)) \geq 2$

<proof>

18 Vertices on Convex Frontier Implies Polygon is Convex

lemma *convex-cut-aux:*

assumes $\forall v \in S. z \cdot v \leq 0$

shows $\text{convex hull } S \subseteq \{x. z \cdot x \leq 0\}$

<proof>

lemma *convex-cut-aux':*

assumes $\forall v \in S. z \cdot v \geq 0$

shows $\text{convex hull } S \subseteq \{x. z \cdot x \geq 0\}$

<proof>

lemma *convex-cut:*

assumes $z \neq 0$

assumes $\{x. z \cdot x = 0\} \cap \text{interior} (\text{convex hull } S) \neq \{\}$

obtains $v1\ v2$ **where** $v1 \neq v2 \wedge \{v1, v2\} \subseteq S \wedge v1 \in \{x. z \cdot x < 0\} \wedge v2 \in \{x. z \cdot x > 0\}$

<proof>

lemma *affine-2-int-convex:*

fixes $S :: 'a::\text{euclidean-space set}$

assumes $\{a, b\} \subseteq S$

assumes $\{a, b\} \subseteq \text{frontier} (\text{convex hull } S)$

assumes $\text{affine hull } \{a, b\} \cap \text{interior} (\text{convex hull } S) \neq \{\}$

shows $\text{affine hull } \{a, b\} \cap \text{convex hull } S = \text{convex hull } \{a, b\}$
(proof)

lemma *halfplane-frontier-affine-hull*:

fixes $b v :: \text{real}^2$
assumes $b \neq 0$
assumes $v \neq 0$
assumes $b \in \{x. v \cdot x = 0\}$
shows $\{x. v \cdot x = 0\} = \text{affine hull } \{0, b\}$
(proof)

lemma *pts-on-convex-frontier-aux*:

assumes *polygon-of p vts*
assumes $vts!0 = 0$
assumes $\text{set } vts \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
shows $\text{path-image } (\text{linepath } (vts!0) (vts!1)) \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
(proof)

lemma *pts-on-convex-frontier-aux'*:

assumes *polygon-of p vts*
assumes $\text{set } vts \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
shows $\text{path-image } (\text{linepath } (vts!0) (vts!1)) \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
(proof)

lemma *pts-on-convex-frontier*:

assumes *polygon-of p vts*
assumes $\text{set } vts \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
assumes $i < \text{length } vts - 1$
shows $\text{path-image } (\text{linepath } (vts!i) (vts!(i+1))) \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
(proof)

lemma *pts-on-frontier-means-path-image-on-frontier*:

assumes *polygon-of p vts*
assumes $\text{set } vts \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
shows $\text{path-image } p \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
(proof)

lemma *pts-on-convex-frontier-interior*:

assumes *polygon-of p vts*
assumes $\text{set } vts \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
shows $\text{path-inside } p = \text{interior } (\text{convex hull } (\text{set } vts))$
(proof)

lemma *pts-subset-frontier*:

assumes *polygon-of p vts*
assumes $\text{set } vts \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
shows $\text{convex } (\text{path-image } p \cup \text{path-inside } p)$
(proof)

lemma *convex-hull-of-nonconvex-polygon-strict-subset-ep*:
assumes *polygon-of p vts*
assumes \neg (*convex (path-image p \cup path-inside p)*)
shows $\{v. v \text{ extreme-point-of } (\text{convex hull } (\text{set } vts))\} \subset \text{set } vts$
 $\langle \text{proof} \rangle$

lemma *convex-hull-of-nonconvex-polygon-strict-subset*:
assumes *polygon-of p vts*
assumes \neg (*convex (path-image p \cup path-inside p)*)
shows $\exists v \in \text{set } vts. v \in \text{interior } (\text{convex hull } (\text{set } vts))$
 $\langle \text{proof} \rangle$

lemma *convex-polygon-means-linepaths-inside*:
fixes $p :: R\text{-to-}R^2$
assumes *polygon-of p vts*
assumes *convex-is: convex hull (set vts) = (path-inside p \cup path-image p)*
assumes *a-in: a \in (path-inside p \cup path-image p)*
assumes *b-in: b \in (path-inside p \cup path-image p)*
shows *path-image (linepath a b) \subseteq (path-inside p \cup path-image p)*
 $\langle \text{proof} \rangle$

end
theory *Polygon-Splitting*
imports
HOL-Analysis.Complete-Measure
Polygon-Jordan-Curve
Polygon-Convex-Lemmas
begin

19 Polygon Splitting

lemma *split-up-a-list-into-3-parts*:
fixes $i j :: \text{nat}$
assumes $i < \text{length } vts \wedge j < \text{length } vts \wedge i < j$
shows
 $vts = (\text{take } i \text{ vts}) @ ((vts ! i) \# ((\text{take } (j - i - 1) (\text{drop } (\text{Suc } i) \text{ vts})) @ (vts ! j) \# \text{drop } (j - i) (\text{drop } (\text{Suc } i) \text{ vts})))$
 $\langle \text{proof} \rangle$

definition *is-polygon-cut* :: $(\text{real}^2) \text{ list} \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{bool}$ **where**
is-polygon-cut vts x y =
 $(x \neq y \wedge$
 $\text{polygon } (\text{make-polygonal-path } vts) \wedge$
 $\{x, y\} \subseteq \text{set } vts \wedge$
 $\text{path-image } (\text{linepath } x \ y) \cap \text{path-image } (\text{make-polygonal-path } vts) = \{x, y\} \wedge$
 $\text{path-image } (\text{linepath } x \ y) \cap \text{path-inside } (\text{make-polygonal-path } vts) \neq \{\})$

definition *is-polygon-cut-path* :: $(\text{real}^2) \text{ list} \Rightarrow R\text{-to-}R^2 \Rightarrow \text{bool}$ **where**


```

is-polygon-cut-path vts cutpath =
  (let x = pathstart cutpath ; y = pathfinish cutpath in
    (x ≠ y ∧
     polygon (make-polygonal-path vts) ∧
     {x, y} ⊆ set vts ∧
     simple-path cutpath ∧
     path-image cutpath ∩ path-image (make-polygonal-path vts) = {x, y} ∧
     path-image cutpath ∩ path-inside (make-polygonal-path vts) ≠ {}))

```

definition *is-polygon-split* ::

```

(real^2) list ⇒ nat ⇒ nat ⇒ bool where
is-polygon-split vts i j =
  (i < length vts ∧ j < length vts ∧ i < j ∧
   (let vts1 = (take i vts) in
    let vts2 = (take (j - i - 1) (drop (Suc i) vts)) in
    let vts3 = drop (j - i) (drop (Suc i) vts) in
    let x = vts ! i in
    let y = vts ! j in
    let p = make-polygonal-path (vts@[vts!0]) in
    let p1 = make-polygonal-path (x#(vts2@[y, x])) in
    let p2 = make-polygonal-path (vts1 @ [x, y] @ vts3 @ [vts ! 0]) in
    let c1 = make-polygonal-path (x#(vts2@[y])) in
    let c2 = make-polygonal-path (vts1 @ [x, y] @ vts3) in
    (is-polygon-cut (vts@[vts!0]) x y ∧
     polygon p ∧ polygon p1 ∧ polygon p2 ∧
     path-inside p1 ∩ path-inside p2 = {} ∧
     path-inside p1 ∪ path-inside p2 ∪ (path-image (linepath x y) - {x, y}) =
     path-inside p
     ∧ ((path-image p1) - (path-image (linepath x y))) ∩ ((path-image p2) -
     (path-image (linepath x y)))
     = {}
     ∧ path-image p
     = ((path-image p1) - (path-image (linepath x y))) ∪ ((path-image p2) -
     (path-image (linepath x y))) ∪ {x, y}
     )))

```

definition *is-polygon-split-path* :: (real^2) list ⇒ nat ⇒ nat ⇒ (real^2) list ⇒ bool **where**

```

is-polygon-split-path vts i j cutvts =
  (i < length vts ∧ j < length vts ∧ i < j ∧
   (let vts1 = (take i vts) in
    let vts2 = (take (j - i - 1) (drop (Suc i) vts)) in
    let vts3 = drop (j - i) (drop (Suc i) vts) in
    let x = vts!i in
    let y = vts!j in
    let cutpath = make-polygonal-path (x # cutvts @ [y]) in
    let p = make-polygonal-path (vts@[vts!0]) in
    let p1 = make-polygonal-path (x#(vts2 @ [y] @ (rev cutvts) @ [x])) in

```

let $p2 = \text{make-polygonal-path } (vts1 \ @ \ ([x] \ @ \ \text{cutvts} \ @ \ [y]) \ @ \ vts3 \ @ \ [vts \ ! \ 0])$ in
 let $c1 = \text{make-polygonal-path } (x\#(vts2@[y]))$ in
 let $c2 = \text{make-polygonal-path } (vts1 \ @ \ ([x] \ @ \ \text{cutvts} \ @ \ [y]) \ @ \ vts3)$ in
 ($\text{is-polygon-cut-path } (vts@[vts!0]) \ \text{cutpath} \ \wedge$
 $\text{polygon } p \ \wedge \ \text{polygon } p1 \ \wedge \ \text{polygon } p2 \ \wedge$
 $\text{path-inside } p1 \ \cap \ \text{path-inside } p2 = \{\}$ \wedge
 $\text{path-inside } p1 \ \cup \ \text{path-inside } p2 \ \cup \ (\text{path-image } \text{cutpath} - \{x, y\}) = \text{path-inside}$
 p
 $\wedge ((\text{path-image } p1) - (\text{path-image } \text{cutpath})) \cap ((\text{path-image } p2) - (\text{path-image}$
 $\text{cutpath})) = \{\}$
 $\wedge \text{path-image } p$
 $= ((\text{path-image } p1) - (\text{path-image } \text{cutpath})) \cup ((\text{path-image } p2) - (\text{path-image}$
 $\text{cutpath})) \cup \{x, y\}$
 $)))$

lemma *polygon-split-add-measure*:

fixes $p \ p1 \ p2 :: R\text{-to-}R2$

assumes *is-polygon-split* $vts \ i \ j$

assumes $vts1 = (\text{take } i \ vts)$

$vts2 = (\text{take } (j - i - 1) \ (\text{drop } (\text{Suc } i) \ vts))$

$vts3 = \text{drop } (j - i) \ (\text{drop } (\text{Suc } i) \ vts)$

$x = vts \ ! \ i$

$y = vts \ ! \ j$

$p = \text{make-polygonal-path } (vts@[vts!0])$

$p1 = \text{make-polygonal-path } (x\#(vts2@[y, x]))$

$p2 = \text{make-polygonal-path } (vts1 \ @ \ [x, y] \ @ \ vts3 \ @ \ [vts \ ! \ 0])$

defines $M1 \equiv \text{measure lebesgue } (\text{path-inside } p1)$ **and**

$M2 \equiv \text{measure lebesgue } (\text{path-inside } p2)$ **and**

$M \equiv \text{measure lebesgue } (\text{path-inside } p)$

shows $M1 + M2 = M$

<proof>

lemma *polygonal-paths-measurable*:

shows $\text{path-image } (\text{make-polygonal-path } vts) \in \text{sets lebesgue}$

<proof>

lemma *polygonal-path-has-emeasure-0*:

shows $\text{emeasure lebesgue } (\text{path-image } (\text{make-polygonal-path } vts)) = 0$

<proof>

lemma *polygon-split-path-add-measure*:

fixes $p \ p1 \ p2 :: R\text{-to-}R2$

assumes *is-polygon-split-path* $vts \ i \ j \ \text{cutvts}$

assumes $vts1 = (\text{take } i \ vts)$

$vts2 = (\text{take } (j - i - 1) \ (\text{drop } (\text{Suc } i) \ vts))$

$vts3 = \text{drop } (j - i) \ (\text{drop } (\text{Suc } i) \ vts)$

$x = vts \ ! \ i$

$y = vts \ ! \ j$

$p = \text{make-polygonal-path } (vts@[vts!0])$

$p1 = \text{make-polygonal-path } (x\#(vts2 \text{ @ } [y] \text{ @ } (\text{rev cutvts}) \text{ @ } [x]))$
 $p2 = \text{make-polygonal-path } (vts1 \text{ @ } ([x] \text{ @ } \text{cutvts} \text{ @ } [y]) \text{ @ } vts3 \text{ @ } [vts ! 0])$
defines $M1 \equiv \text{measure lebesgue } (\text{path-inside } p1)$ **and**
 $M2 \equiv \text{measure lebesgue } (\text{path-inside } p2)$ **and**
 $M \equiv \text{measure lebesgue } (\text{path-inside } p)$
shows $M1 + M2 = M$
 $\langle \text{proof} \rangle$

lemma *polygon-cut-path-to-split-path-vtx0*:
fixes $p :: R\text{-to-}R2$
assumes *polygon-p*: *polygon p* **and**
 $i\text{-gt}: i > 0$ **and**
 $i\text{-lt}: i < \text{length } vts$ **and**
 $p\text{-is}: p = \text{make-polygonal-path } (vts \text{ @ } [vts ! 0])$ **and**
 $\text{cutpath}: \text{cutpath} = \text{make-polygonal-path } ([vts!0] \text{ @ } \text{cutvts} \text{ @ } [vts!i])$ **and**
 $\text{have-cut}: \text{is-polygon-cut-path } (vts \text{ @ } [vts!0]) \text{ cutpath}$
shows *is-polygon-split-path vts 0 i cutvts*
 $\langle \text{proof} \rangle$

lemma *polygon-cut-path-to-split-path*:
fixes $p :: R\text{-to-}R2$
assumes *polygon p*
 $p = \text{make-polygonal-path } (vts \text{ @ } [vts ! 0])$
 $\text{is-polygon-cut-path } (vts \text{ @ } [vts!0]) \text{ cutpath}$
 $vts1 \equiv (\text{take } i \text{ } vts)$
 $vts2 \equiv (\text{take } (j - i - 1) (\text{drop } (\text{Suc } i) \text{ } vts))$
 $vts3 \equiv \text{drop } (j - i) (\text{drop } (\text{Suc } i) \text{ } vts)$
 $x \equiv vts ! i$
 $y \equiv vts ! j$
 $\text{cutpath} = \text{make-polygonal-path } ([x] \text{ @ } \text{cutvts} \text{ @ } [y])$
 $i < \text{length } vts \wedge j < \text{length } vts \wedge i < j$
 $p1 \equiv \text{make-polygonal-path } (x\#(vts2\text{@}([y] \text{ @ } (\text{rev cutvts}) \text{ @ } [x])))$ **and**
 $p2 \equiv \text{make-polygonal-path } (vts1 \text{ @ } ([x] \text{ @ } \text{cutvts} \text{ @ } [y]) \text{ @ } vts3 \text{ @ } [(vts1 \text{ @ } [x]) ! 0])$
shows *is-polygon-split-path vts i j cutvts*
 $\langle \text{proof} \rangle$

lemma *good-polygonal-path-implies-polygon-split-path*:
assumes *polygon p*
assumes $p = \text{make-polygonal-path } (vts \text{ @ } [vts!0])$
assumes *good-polygonal-path v1 cutvts v2* $(vts \text{ @ } [vts!0])$
assumes $i < \text{length } vts \wedge j < \text{length } vts$
assumes $vts ! i = v1$
assumes $vts ! j = v2$
assumes $i < j$
shows *is-polygon-split-path vts i j cutvts*
 $\langle \text{proof} \rangle$

```

lemma good-path-iff:
  good-linepath a b vts  $\longleftrightarrow$  good-polygonal-path a [] b vts
  <proof>

lemma polygon-cut-iff: is-polygon-cut (vts @ [vts!0]) (vts!i) (vts!j)
   $\longleftrightarrow$  is-polygon-cut-path (vts @ [vts!0]) (linepath (vts!i) (vts!j))
  <proof>

lemma polygon-split-iff: is-polygon-split vts i j  $\longleftrightarrow$  is-polygon-split-path vts i j []
  <proof>

lemma polygon-cut-to-split-vtx0:
  fixes p :: R-to-R2
  assumes polygon-p: polygon p and
    i-gt: i > 0 and
    i-lt: i < length vts and
    p-is: p = make-polygonal-path (vts @ [vts ! 0]) and
    have-cut: is-polygon-cut (vts @ [vts!0]) (vts!0) (vts!i)
  shows is-polygon-split vts 0 i
  <proof>

lemma polygon-cut-to-split:
  fixes p :: R-to-R2
  assumes is-polygon-cut (vts @ [vts!0]) (vts!i) (vts!j)
    i < length vts  $\wedge$  j < length vts  $\wedge$  i < j
  shows is-polygon-split vts i j
  <proof>

lemma good-linepath-implies-polygon-split:
  assumes polygon p
  assumes p = make-polygonal-path (vts @ [vts!0])
  assumes good-linepath v1 v2 (vts @ [vts!0])
  assumes i < length vts  $\wedge$  j < length vts
  assumes vts ! i = v1
  assumes vts ! j = v2
  assumes i < j
  shows is-polygon-split vts i j
  <proof>

end
theory Triangle-Lemmas
imports
  Polygon-Convex-Lemmas
  Integral-Matrix
  Affine-Arithmetic.Floatarith-Expression
  HOL-Analysis.Topology-Euclidean-Space
  HOL-Analysis.Equivalence-Lebesgue-Henstock-Integration
  HOL-Analysis.Inner-Product

```

HOL-Analysis.Line-Segment
HOL-Analysis.Convex-Euclidean-Space
HOL-Analysis.Change-Of-Vars

begin

20 Triangles

definition *elem-triangle* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{bool}$ **where**

elem-triangle $a\ b\ c \iff$
 $\neg \text{collinear}\ \{a,\ b,\ c\}$
 $\wedge \text{integral-vec}\ a \wedge \text{integral-vec}\ b \wedge \text{integral-vec}\ c$
 $\wedge \{x.\ x \in \text{convex hull}\ \{a,\ b,\ c\} \wedge \text{integral-vec}\ x\} = \{a,\ b,\ c\}$

definition *triangle-mat* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^{2 \times 2}$ **where**

triangle-mat $a\ b\ c = \text{transpose}\ (\text{vector}\ [b - a,\ c - a])$

definition *triangle-linear* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow (\text{real}^2 \Rightarrow \text{real}^2)$
where

triangle-linear $a\ b\ c = (\lambda x.\ (\text{triangle-mat}\ a\ b\ c) *v\ x)$

definition *triangle-affine* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow (\text{real}^2 \Rightarrow \text{real}^2)$ **where**

triangle-affine $a\ b\ c = (\lambda x.\ a + (\text{triangle-mat}\ a\ b\ c) *v\ x)$

abbreviation *unit-square* \equiv

$(\text{convex hull}\ \{\text{vector}\ [0,\ 0],\ \text{vector}\ [0,\ 1],\ \text{vector}\ [1,\ 1],\ \text{vector}\ [1,\ 0]\})::(\text{real}^2)$
set)

abbreviation *unit-triangle* \equiv

$(\text{convex hull}\ \{\text{vector}\ [0,\ 0],\ \text{vector}\ [1,\ 0],\ \text{vector}\ [0,\ 1]\})::(\text{real}^2)$ *set*)

abbreviation *unit-triangle'* \equiv

$(\text{convex hull}\ \{\text{vector}\ [1,\ 1],\ \text{vector}\ [1,\ 0],\ \text{vector}\ [0,\ 1]\})::(\text{real}^2)$ *set*)

lemma *triangle-inside-is-convex-hull-interior*:

assumes *polygon-of* $p\ [a,\ b,\ c,\ a]$

shows *path-inside* $p = \text{interior}\ (\text{convex hull}\ \{a,\ b,\ c\})$

<proof>

lemma *triangle-is-convex*:

assumes $p = \text{make-triangle}\ a\ b\ c$ **and** $\neg \text{collinear}\ \{a,\ b,\ c\}$

shows *convex* (*path-inside* p) (**is convex** ?s)

<proof>

lemma *affine-comp-linear-trans*: *triangle-affine* $a\ b\ c = (\lambda x.\ x + a) \circ (\text{triangle-linear}\ a\ b\ c)$

<proof>

lemma *triangle-linear-der*:

fixes $a\ b\ c :: \text{real}^2$
defines $T \equiv \text{triangle-linear } a\ b\ c$
shows $(T \text{ has-derivative } T)$ (at x)
 <proof>

lemma *triangle-affine-der*:
fixes $a\ b\ c :: \text{real}^2$
assumes $S \in \text{sets lebesgue}$ **and** $x \in S$
defines $A \equiv \text{triangle-affine } a\ b\ c$ **and** $T \equiv \text{triangle-linear } a\ b\ c$
shows $x \in S \implies (A \text{ has-derivative } T)$ (at x within S)
 <proof>

lemma *triangle-linear-inj*:
fixes $a\ b\ c :: \text{real}^2$
assumes $\neg \text{collinear } \{a, b, c\}$
defines $L \equiv \text{triangle-linear } a\ b\ c$
shows *inj* L
 <proof>

lemma *triangle-affine-inj*:
fixes $a\ b\ c :: \text{real}^2$
assumes $\neg \text{collinear } \{a, b, c\}$
defines $A \equiv \text{triangle-affine } a\ b\ c$
shows *inj* A
 <proof>

lemma *triangle-linear-integrable*:
fixes $a\ b\ c :: \text{real}^2$
assumes $S \in \text{lmeasurable}$
defines $T \equiv \text{triangle-linear } a\ b\ c$
shows $(\lambda x. \text{abs } (\det (\text{matrix } (T))))$ *integrable-on* S (**is** $(\lambda x. ?c)$ *integrable-on* S)
 <proof>

lemma *measure-differentiable-image-eq-affine*:
fixes $a\ b\ c :: \text{real}^2$
defines $A \equiv \text{triangle-affine } a\ b\ c$ **and** $T \equiv \text{triangle-linear } a\ b\ c$
assumes $S \in \text{lmeasurable}$ **and** $\neg \text{collinear } \{a, b, c\}$
shows $\text{measure lebesgue } (A \text{ ' } S) = \text{integral } S (\lambda x. \text{abs } (\det (\text{matrix } T)))$
 <proof>

lemma *triangle-affine-img*:
fixes $a\ b\ c :: \text{real}^2$
defines $A \equiv \text{triangle-affine } a\ b\ c$
shows $\text{convex hull } \{a, b, c\} = A \text{ ' unit-triangle}$
 <proof>

lemma *triangle-affine-e1-e2*:
fixes $a\ b\ c :: \text{real}^2$
defines $A \equiv \text{triangle-affine } a\ b\ c$

shows (*triangle-affine* $a\ b\ c$) (*vector* $[0, 0]$) = a
 (*triangle-affine* $a\ b\ c$) (*vector* $[1, 0]$) = b
 (*triangle-affine* $a\ b\ c$) (*vector* $[0, 1]$) = c
 ⟨*proof*⟩

lemma *triangle-measure-integral-of-det*:

fixes $a\ b\ c :: \text{real}^2$
defines $S \equiv \text{convex hull } \{a, b, c\}$
assumes $\neg \text{collinear } \{a, b, c\}$
shows $\text{measure lebesgue } S =$
 $\text{integral unit-triangle } (\lambda(x::\text{real}^2). \text{abs } (\text{det } (\text{matrix } (\text{triangle-linear } a\ b$
 $c))))$
 ⟨*proof*⟩

lemma *triangle-affine-preserves-interior*:

assumes $A = \text{triangle-affine } a\ b\ c$ **and** $L = \text{triangle-linear } a\ b\ c$
assumes $\neg \text{collinear } \{a, b, c\}$
shows $A \text{ ' } (\text{interior } S) = \text{interior } (A \text{ ' } S)$
 ⟨*proof*⟩

lemma *triangle-affine-preserves-affine-hull*:

assumes $A = \text{triangle-affine } a\ b\ c$
assumes $\neg \text{collinear } \{a, b, c\}$
shows $A \text{ ' } (\text{affine hull } S) = \text{affine hull } (A \text{ ' } S)$
 ⟨*proof*⟩

lemma *triangle-measure-convex-hull-measure-path-inside-same*:

assumes $p\text{-triangle}: p = \text{make-triangle } a\ b\ c$
assumes $\text{elem-triangle}: \text{elem-triangle } a\ b\ c$
shows $\text{measure lebesgue } (\text{convex hull } \{a, b, c\}) = \text{measure lebesgue } (\text{path-inside } p)$
 $(\text{is measure lebesgue } ?S = \text{measure lebesgue } ?I)$
 ⟨*proof*⟩

lemma *on-triangle-path-image-cases*:

assumes $p = \text{make-triangle } a\ b\ c$
assumes $d \in \text{path-image } p$
shows $d \in \text{path-image } (\text{linepath } a\ b) \vee d \in \text{path-image } (\text{linepath } b\ c) \vee d \in$
 $\text{path-image } (\text{linepath } c\ a)$
 ⟨*proof*⟩

lemma *on-triangle-frontier-cases*:

fixes $a\ b\ c :: \text{real}^2$
assumes $\neg \text{collinear } \{a, b, c\}$
assumes $d \in \text{frontier } (\text{convex hull } \{a, b, c\})$
shows $d \in \text{path-image } (\text{linepath } a\ b) \vee d \in \text{path-image } (\text{linepath } b\ c) \vee d \in$
 $\text{path-image } (\text{linepath } c\ a)$
 ⟨*proof*⟩

lemma *triangle-path-image-subset-convex*:
assumes $p = \text{make-triangle } a \ b \ c$
shows $\text{path-image } p \subseteq \text{convex hull } \{a, b, c\}$
 $\langle \text{proof} \rangle$

lemma *triangle-convex-hull*:
assumes $p = \text{make-triangle } a \ b \ c$ **and** $\neg \text{collinear } \{a, b, c\}$
shows $\text{convex hull } \{a, b, c\} = (\text{path-image } p) \cup (\text{path-inside } p)$
 $\langle \text{proof} \rangle$

end
theory *Unit-Geometry*
imports
HOL-Analysis.Polytope
Polygon-Jordan-Curve
Triangle-Lemmas

begin

21 Measure Setup

lemma *finite-convex-is-measurable*:
fixes $p :: (\text{real}^2) \text{ set}$
assumes $p = \text{convex hull } l$ **and** *finite* l
shows $p \in \text{sets lebesgue}$
 $\langle \text{proof} \rangle$

lemma *unit-square-lebesgue*: $\text{unit-square} \in \text{sets lebesgue}$
 $\langle \text{proof} \rangle$

lemma *unit-triangle-lebesgue*: $\text{unit-triangle} \in \text{sets lebesgue}$
 $\langle \text{proof} \rangle$

lemma *unit-triangle-lmeasurable*: $\text{unit-triangle} \in \text{lmeasurable}$
 $\langle \text{proof} \rangle$

22 Unit Triangle

lemma *unit-triangle-vts-not-collinear*:
 $\neg \text{collinear } \{(\text{vector } [0, 0])::\text{real}^2, \text{vector } [1, 0], \text{vector } [0, 1]\}$
(is $\neg \text{collinear } \{?a, ?b, ?c\}$ **)**
 $\langle \text{proof} \rangle$

lemma *unit-triangle-convex*:
assumes $p = (\text{make-polygonal-path } [\text{vector } [0, 0], \text{vector } [1, 0], \text{vector } [0, 1], \text{vector } [0, 0]])$
(is $p = \text{make-polygonal-path } [?O, ?e1, ?e2, ?O]$ **)**

shows *convex* (*path-inside p*)
 ⟨*proof*⟩

lemma *unit-triangle-char*:
shows *unit-triangle* = {*x*. $0 \leq x \$ 1 \wedge 0 \leq x \$ 2 \wedge x \$ 1 + x \$ 2 \leq 1$ }
 (**is** *unit-triangle* = ?*S*)
 ⟨*proof*⟩

lemma *unit-triangle-interior-char*:
shows *interior unit-triangle* = {*x*. $0 < x \$ 1 \wedge 0 < x \$ 2 \wedge x \$ 1 + x \$ 2 < 1$ }
 (**is** *interior unit-triangle* = ?*S*)
 ⟨*proof*⟩

lemma *unit-triangle-is-elementary*: *elem-triangle* (*vector* [0, 0]) (*vector* [1, 0])
 (*vector* [0, 1])
 (**is** *elem-triangle* ?*a* ?*b* ?*c*)
 ⟨*proof*⟩

lemma *unit-triangles-same-area*:
measure lebesgue unit-triangle' = *measure lebesgue unit-triangle*
 ⟨*proof*⟩

23 Unit Square

lemma *convex-hull-4*:
convex hull {*a,b,c,d*} = { *u *_R a + v *_R b + w *_R c + t *_R d* | *u v w t*. $0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge 0 \leq t \wedge u + v + w + t = 1$ }
 ⟨*proof*⟩

lemma *unit-square-characterization-helper*:
fixes *a b* :: *real*
assumes $0 \leq a \wedge a \leq 1 \wedge 0 \leq b \wedge b \leq 1$ **and**
 $a \leq b$
obtains *u v w t* **where**
 $\text{vector } [a, b] = u *_R ((\text{vector } [0, 0]))::\text{real}^2$
 $+ v *_R (\text{vector } [0, 1])$
 $+ w *_R (\text{vector } [1, 1])$
 $+ t *_R (\text{vector } [1, 0])$
 $\wedge 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge 0 \leq t \wedge u + v + w + t = 1$
 ⟨*proof*⟩

lemma *unit-square-characterization*:
unit-square = {*x*. $0 \leq x \$ 1 \wedge x \$ 1 \leq 1 \wedge 0 \leq x \$ 2 \wedge x \$ 2 \leq 1$ } (**is** *unit-square*
 = ?*S*)
 ⟨*proof*⟩

lemma *e1e2-basis*:
defines *e1* ≡ (*vector* [1, 0])::(*real*²) **and**

$e2 \equiv (\text{vector } [0, 1]) :: (\text{real}^2)$
shows $e1 = \text{axis } 1 \ (1 :: \text{real})$ **and** $e1 \in (\text{Basis} :: ((\text{real}^2) \text{ set}))$ **and**
 $e2 = \text{axis } 2 \ (1 :: \text{real})$ **and** $e2 \in (\text{Basis} :: ((\text{real}^2) \text{ set}))$
 <proof>

lemma *unit-square-cbox*: $\text{unit-square} = \text{cbox } (\text{vector } [0, 0]) \ (\text{vector } [1, 1])$
 <proof>

lemma *unit-square-area*: $\text{measure lebesgue unit-square} = 1$
 <proof>

24 Unit Triangle Area is 1/2

lemma *unit-triangle'-char*:
shows $\text{unit-triangle}' = \{x. x \$ 1 \leq 1 \wedge x \$ 2 \leq 1 \wedge x \$ 1 + x \$ 2 \geq 1\}$
 <proof>

lemma *unit-square-split-diag*:
shows $\text{unit-square} = \text{unit-triangle} \cup \text{unit-triangle}'$
 <proof>

lemma *unit-triangle-INT-unit-triangle'-measure*:
 $\text{measure lebesgue } (\text{unit-triangle} \cap \text{unit-triangle}') = 0$
 <proof>

lemma *unit-triangle-area*: $\text{measure lebesgue unit-triangle} = 1/2$
 <proof>

end
theory *Elementary-Triangle-Area*
imports
Unit-Geometry

begin

25 Area of Elementary Triangle is 1/2

lemma *nonint-in-square-img-IMP-nonint-triangle-img*:
assumes $A = \text{triangle-affine } a \ b \ c$
assumes $x \in \text{unit-square}$
assumes $\neg \text{integral-vec } x$
assumes $\text{integral-vec } (A \ x)$
assumes $\text{elem-triangle } a \ b \ c$
obtains x' **where** $x' \in \text{unit-triangle} \wedge \neg \text{integral-vec } x' \wedge \text{integral-vec } (A \ x')$
 <proof>

lemma *elem-triangle-integral-mat-bij*:
fixes $a \ b \ c :: \text{real}^2$

```

assumes elem-triangle a b c
defines  $L \equiv \text{triangle-mat } a \ b \ c$ 
shows integral-mat-bij  $L$ 
<proof>

lemma elem-triangle-measure-integral-of-1:
fixes  $a \ b \ c :: \text{real}^2$ 
defines  $S \equiv \text{convex hull } \{a, b, c\}$ 
assumes elem-triangle  $a \ b \ c$ 
shows measure lebesgue  $S = \text{integral unit-triangle } (\lambda(x::\text{real}^2). 1)$ 
<proof>

lemma elem-triangle-area-is-half:
fixes  $a \ b \ c :: \text{real}^2$ 
assumes elem-triangle  $a \ b \ c$ 
defines  $S \equiv \text{convex hull } \{a, b, c\}$ 
shows measure lebesgue  $S = 1/2$  (is  $?S\text{-area} = 1/2$ )
<proof>

end
theory Pick
imports
  Polygon-Splitting
  Elementary-Triangle-Area
begin

```

26 Setup

26.1 Integral Points Cardinality Properties

```

lemma bounded-finite:
fixes  $A::(\text{real}^2)$  set
assumes bounded  $A$ 
shows finite  $\{x::(\text{real}^2). \text{integral-vec } x \wedge x \in A\}$  (is finite  $?A\text{-int}$ )
<proof>

```

```

lemma finite-path-image:
assumes polygon  $p$ 
shows finite  $\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
<proof>

```

```

lemma finite-path-inside:
assumes polygon  $p$ 
shows finite  $\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$ 
<proof>

```

```

lemma bounded-finite-inside:
fixes  $B::(\text{real}^2)$  set
assumes simple-path  $p$ 

```

shows *bounded* (*path-inside* *p*)
 ⟨*proof*⟩

lemma *finite-integral-points-path-image*:
assumes *simple-path* *p*
shows *finite* {*x*. *integral-vec* *x* ∧ *x* ∈ *path-image* *p*}
 ⟨*proof*⟩

lemma *finite-integral-points-path-inside*:
assumes *simple-path* *p*
shows *finite* {*x*. *integral-vec* *x* ∧ *x* ∈ *path-inside* *p*}
 ⟨*proof*⟩

27 Pick splitting

lemma *pick-split-path-union-main*:
assumes *is-split*: *is-polygon-split-path* *vts* *i* *j* *cutvts*
assumes *vts1* = (*take* *i* *vts*)
assumes *vts2* = (*take* (*j* - *i* - 1) (*drop* (*Suc* *i*) *vts*))
assumes *vts3* = *drop* (*j* - *i*) (*drop* (*Suc* *i*) *vts*)
assumes *x* = *vts!**i*
assumes *y* = *vts!**j*
assumes *cutpath* = *make-polygonal-path* (*x* # *cutvts* @ [*y*])
assumes *p*: *p* = *make-polygonal-path* (*vts*@[*vts!*0]) (**is** *p* = *make-polygonal-path* ?*p-vts*)
assumes *p1*: *p1* = *make-polygonal-path* (*x*#(*vts2* @ [*y*] @ (*rev* *cutvts*) @ [*x*]))
(is *p1* = *make-polygonal-path* ?*p1-vts*)
assumes *p2*: *p2* = *make-polygonal-path* (*vts1* @ ([*x*] @ *cutvts* @ [*y*]) @ *vts3* @ [*vts* ! 0]) (**is** *p2* = *make-polygonal-path* ?*p2-vts*)
assumes *I1*: *I1* = *card* {*x*. *integral-vec* *x* ∧ *x* ∈ *path-inside* *p1*}
assumes *B1*: *B1* = *card* {*x*. *integral-vec* *x* ∧ *x* ∈ *path-image* *p1*}
assumes *I2*: *I2* = *card* {*x*. *integral-vec* *x* ∧ *x* ∈ *path-inside* *p2*}
assumes *B2*: *B2* = *card* {*x*. *integral-vec* *x* ∧ *x* ∈ *path-image* *p2*}
assumes *I*: *I* = *card* {*x*. *integral-vec* *x* ∧ *x* ∈ *path-inside* *p*}
assumes *B*: *B* = *card* {*x*. *integral-vec* *x* ∧ *x* ∈ *path-image* *p*}
assumes *all-integral-vts*: *all-integral* *vts*
shows *measure lebesgue* (*path-inside* *p1*) = *I1* + *B1*/2 - 1
 ⇒ *measure lebesgue* (*path-inside* *p2*) = *I2* + *B2*/2 - 1
 ⇒ *measure lebesgue* (*path-inside* *p*) = *I* + *B*/2 - 1
measure lebesgue (*path-inside* *p*) = *I* + *B*/2 - 1
 ⇒ *measure lebesgue* (*path-inside* *p2*) = *I2* + *B2*/2 - 1
 ⇒ *measure lebesgue* (*path-inside* *p1*) = *I1* + *B1*/2 - 1
measure lebesgue (*path-inside* *p*) = *I* + *B*/2 - 1
 ⇒ *measure lebesgue* (*path-inside* *p1*) = *I1* + *B1*/2 - 1
 ⇒ *measure lebesgue* (*path-inside* *p2*) = *I2* + *B2*/2 - 1
 ⟨*proof*⟩

lemma *pick-split-union*:
assumes *is-split*: *is-polygon-split* *vts* *i* *j*

assumes $vts1 = (take\ i\ vts)$
assumes $vts2 = (take\ (j - i - 1)\ (drop\ (Suc\ i)\ vts))$
assumes $vts3 = drop\ (j - i)\ (drop\ (Suc\ i)\ vts)$
assumes $x = vts!\ i$
assumes $y = vts!\ j$
assumes $p: p = make\text{-}polygonal\text{-}path\ (vts@[vts!0])$ (**is** $p = make\text{-}polygonal\text{-}path\ ?p\text{-}vts$)
assumes $p1: p1 = make\text{-}polygonal\text{-}path\ (x\#\ (vts2@[y, x]))$ (**is** $p1 = make\text{-}polygonal\text{-}path\ ?p1\text{-}vts$)
assumes $p2: p2 = make\text{-}polygonal\text{-}path\ (vts1\ @\ [x, y]\ @\ vts3\ @\ [vts!\ 0])$ (**is** $p2 = make\text{-}polygonal\text{-}path\ ?p2\text{-}vts$)
assumes $I1: I1 = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}inside\ p1\}$
assumes $B1: B1 = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}image\ p1\}$
assumes $pick1: measure\ lebesgue\ (path\text{-}inside\ p1) = I1 + B1/2 - 1$
assumes $I2: I2 = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}inside\ p2\}$
assumes $B2: B2 = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}image\ p2\}$
assumes $pick2: measure\ lebesgue\ (path\text{-}inside\ p2) = I2 + B2/2 - 1$
assumes $I: I = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}inside\ p\}$
assumes $B: B = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}image\ p\}$
assumes $all\text{-}integral\text{-}vts: all\text{-}integral\ vts$
shows $measure\ lebesgue\ (path\text{-}inside\ p) = I + B/2 - 1$
 $measure\ lebesgue\ (path\text{-}inside\ p) = measure\ lebesgue\ (path\text{-}inside\ p1) +$
 $measure\ lebesgue\ (path\text{-}inside\ p2)$
 $\langle proof \rangle$

lemma *pick-split-path-union:*

assumes *is-split: is-polygon-split-path* $vts\ i\ j\ cutvts$
assumes $vts1 = (take\ i\ vts)$
assumes $vts2 = (take\ (j - i - 1)\ (drop\ (Suc\ i)\ vts))$
assumes $vts3 = drop\ (j - i)\ (drop\ (Suc\ i)\ vts)$
assumes $x = vts!\ i$
assumes $y = vts!\ j$
assumes $cutpath = make\text{-}polygonal\text{-}path\ (x\ #\ cutvts\ @\ [y])$
assumes $p: p = make\text{-}polygonal\text{-}path\ (vts@[vts!0])$ (**is** $p = make\text{-}polygonal\text{-}path\ ?p\text{-}vts$)
assumes $p1: p1 = make\text{-}polygonal\text{-}path\ (x\#\ (vts2\ @\ [y]\ @\ (rev\ cutvts)\ @\ [x]))$
(is $p1 = make\text{-}polygonal\text{-}path\ ?p1\text{-}vts$)
assumes $p2: p2 = make\text{-}polygonal\text{-}path\ (vts1\ @\ ([x]\ @\ cutvts\ @\ [y])\ @\ vts3\ @\ [vts!\ 0])$ (**is** $p2 = make\text{-}polygonal\text{-}path\ ?p2\text{-}vts$)
assumes $I1: I1 = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}inside\ p1\}$
assumes $B1: B1 = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}image\ p1\}$
assumes $pick1: measure\ lebesgue\ (path\text{-}inside\ p1) = I1 + B1/2 - 1$
assumes $I2: I2 = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}inside\ p2\}$
assumes $B2: B2 = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}image\ p2\}$
assumes $pick2: measure\ lebesgue\ (path\text{-}inside\ p2) = I2 + B2/2 - 1$
assumes $I: I = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}inside\ p\}$
assumes $B: B = card\ \{x.\ integral\text{-}vec\ x \wedge x \in path\text{-}image\ p\}$
assumes $all\text{-}integral\text{-}vts: all\text{-}integral\ vts$
shows $measure\ lebesgue\ (path\text{-}inside\ p) = I + B/2 - 1$

<proof>

lemma *pick-triangle-basic-split*:

assumes $p = \text{make-triangle } a \ b \ c$ **and** $\text{distinct } [a, b, c]$ **and** $\neg \text{collinear } \{a, b, c\}$ **and**

d-prop: $d \in \text{path-image } (\text{linepath } a \ b) \wedge d \notin \{a, b, c\}$

shows $\text{good-linepath } c \ d \ [a, d, b, c, a]$

$\wedge \text{path-image } (\text{make-polygonal-path } [a, d, b, c, a]) = \text{path-image } p$

<proof>

28 Convex Hull Has Good Linepath

lemma *leq-2-extreme-points-means-collinear*:

fixes $pts :: 'a::\text{euclidean-space set}$

assumes $\text{finite } pts$

assumes $\text{card } \{v. v \ \text{extreme-point-of } (\text{convex hull } pts)\} \leq 2$

shows $\text{collinear } pts$

<proof>

lemma *convex-hull-non-extreme-point-in-open-seg*:

assumes $H = \text{convex hull } pts$

assumes $x \in H - \{v. v \ \text{extreme-point-of } H\}$

shows $\exists a \ b. a \in H \wedge b \in H \wedge x \in \text{open-segment } a \ b$

<proof>

lemma *convex-hull-extreme-points-vertex-split*:

fixes $pts :: (\text{real}^2) \ \text{set}$

assumes $H = \text{convex hull } pts$

assumes $\text{finite } pts$

assumes $\text{card } \{v. v \ \text{extreme-point-of } H\} \geq 4$

assumes $\{a, b, c\} \subseteq \{v. v \ \text{extreme-point-of } H\} \wedge \text{distinct } [a, b, c]$

shows $\text{path-image } (\text{linepath } a \ b) \cap \text{interior } H \neq \{\}$

$\vee \text{path-image } (\text{linepath } b \ c) \cap \text{interior } H \neq \{\}$

$\vee \text{path-image } (\text{linepath } c \ a) \cap \text{interior } H \neq \{\}$

<proof>

lemma *convex-hull-has-vertex-split-helper-wlog*:

assumes $p = \text{make-triangle } a \ b \ c$ **and** $\text{distinct } [a, b, c]$ **and** $\neg \text{collinear } \{a, b, c\}$ **and**

d-prop: $d \in \text{path-image } (\text{linepath } a \ b) \wedge d \notin \{a, b, c\}$

shows $\text{path-image } (\text{linepath } c \ d) \cap \text{path-inside } p \neq \{\}$

<proof>

lemma *convex-hull-has-vertex-split-helper*:

assumes $p = \text{make-triangle } a \ b \ c$ **and** $\text{distinct } [a, b, c]$ **and** $\neg \text{collinear } \{a, b, c\}$ **and**

d-prop: $d \in \text{path-image } p \wedge d \notin \{a, b, c\}$

shows $\exists x \ y. \{x, y\} \subseteq \{a, b, c, d\} \wedge x \neq y \wedge \text{path-image } (\text{linepath } x \ y) \cap \text{path-inside } p \neq \{\}$

<proof>

lemma *convex-hull-has-vertex-split:*

fixes *vts* :: (real²) set

assumes $H = \text{convex hull } vts$

assumes $\neg \text{collinear } vts$

assumes $\text{card } vts > 3$

assumes *finite vts*

shows $\exists a b. \{a, b\} \subseteq vts \wedge a \neq b \wedge \text{path-image } (\text{linepath } a \ b) \cap \text{interior } H \neq$

$\{\}$

<proof>

lemma *convex-polygon-has-good-linepath-helper:*

assumes *polygon-of p vts*

assumes $\text{convex } (\text{path-inside } p \cup \text{path-image } p)$

assumes $\text{card } (\text{set } vts) > 3$

obtains *a b* **where** $\{a, b\} \subseteq \text{set } vts \wedge a \neq b \wedge \neg \text{path-image } (\text{linepath } a \ b) \subseteq$

$\text{path-image } p$

<proof>

lemma *convex-polygon-has-good-linepath:*

assumes $\text{convex } (\text{path-inside } p \cup \text{path-image } p)$

assumes *polygon p*

assumes $p = \text{make-polygonal-path } vts$

assumes $\text{card } (\text{set } vts) > 3$

shows $\exists a b. \text{good-linepath } a \ b \ vts$

<proof>

29 Pick's Theorem

definition *integral-inside:*

$\text{integral-inside } p = \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$

definition *integral-boundary:*

$\text{integral-boundary } p = \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$

29.1 Pick's Theorem Triangle Case

definition *pick-triangle:*

$\text{pick-triangle } p \ a \ b \ c \longleftrightarrow$

$p = \text{make-triangle } a \ b \ c$

$\wedge \text{all-integral } [a, b, c]$

$\wedge \text{distinct } [a, b, c]$

$\wedge \neg \text{collinear } \{a, b, c\}$

definition *pick-holds:*

$\text{pick-holds } p \longleftrightarrow$

(let $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$ in

let $B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ in

$$\text{measure lebesgue } (\text{path-inside } p) = I + B/2 - 1$$

lemma *pick-triangle-wlog-helper*:

assumes *pick-triangle* p a b c **and**

$I = \text{card } (\text{integral-inside } p)$ **and**

$B = \text{card } (\text{integral-boundary } p)$ **and**

$\text{integral-inside } p = \{\}$ **and**

$\text{integral-vec } d \wedge d \in \text{path-image } (\text{linepath } a \ b) \wedge d \notin \{a, b, c\}$ **and** $d \notin \{a, b, c\}$ **and**

ih: $\bigwedge p' \ a' \ b' \ c'. (\text{card } (\text{integral-inside } p') + \text{card } (\text{integral-boundary } p') < I + B) \implies \text{pick-triangle } p' \ a' \ b' \ c' \implies \text{pick-holds } p'$

shows $\text{measure lebesgue } (\text{path-inside } p) = I + B/2 - 1$

<proof>

lemma *pick-triangle-helper*:

assumes *pick-triangle* p a b c **and**

$I = \text{card } (\text{integral-inside } p)$ **and**

$B = \text{card } (\text{integral-boundary } p)$ **and**

$\text{integral-inside } p = \{\}$ **and**

$\text{integral-vec } d \wedge d \notin \{a, b, c\}$ **and** $d \notin \{a, b, c\}$ **and**

$d \in \text{path-image } (\text{linepath } a \ b)$

$\vee d \in \text{path-image } (\text{linepath } b \ c)$

$\vee d \in \text{path-image } (\text{linepath } c \ a)$ **and**

ih: $\bigwedge p' \ a' \ b' \ c'. (\text{card } (\text{integral-inside } p') + \text{card } (\text{integral-boundary } p') < I + B) \implies \text{pick-triangle } p' \ a' \ b' \ c' \implies \text{pick-holds } p'$

shows $\text{measure lebesgue } (\text{path-inside } p) = I + B/2 - 1$

<proof>

lemma *triangle-3-split-helper*:

fixes $a \ b :: 'a :: \text{euclidean-space}$

assumes $a \in \text{frontier } S$

assumes $b \in \text{interior } S$

assumes *convex* S

assumes *closed* S

shows $\text{path-image } (\text{linepath } a \ b) \cap \text{frontier } S = \{a\}$

<proof>

lemma *unit-triangle-interior-point-not-collinear-e1-e2*:

assumes $p = \text{make-triangle } (\text{vector } [0, 0]) (\text{vector } [1, 0]) (\text{vector } [0, 1])$

(**is** $p = \text{make-triangle } ?O \ ?e1 \ ?e2$)

assumes $z \in \text{path-inside } p$

shows $\neg \text{collinear } \{?O, ?e1, z\}$

<proof>

lemma *triangle-interior-point-not-collinear-vertices-wlog-helper*:

assumes $p = \text{make-triangle } a \ b \ c$

assumes *polygon* p

assumes $z \in \text{path-inside } p$

shows $\neg \text{collinear } \{a, b, z\}$

<proof>

lemma *triangle-interior-point-not-collinear-vertices:*

assumes $p = \text{make-triangle } a \ b \ c$

assumes *polygon* p

assumes $z \in \text{path-inside } p$

shows $\neg \text{collinear } \{a, b, z\} \wedge \neg \text{collinear } \{a, c, z\} \wedge \neg \text{collinear } \{b, c, z\}$

<proof>

lemma *triangle-3-split:*

assumes $p = \text{make-triangle } a \ b \ c$

assumes *polygon* p

assumes $z \in \text{path-inside } p$

shows *is-polygon-split-path* $[a, b, c] \ 0 \ 1 \ [z]$

is-polygon-split $[a, z, b, c] \ 1 \ 3$

$a \notin \text{path-image } (\text{make-triangle } z \ b \ c) \cup \text{path-inside } (\text{make-triangle } z \ b \ c)$

$b \notin \text{path-image } (\text{make-triangle } a \ z \ c) \cup \text{path-inside } (\text{make-triangle } a \ z \ c)$

$c \notin \text{path-image } (\text{make-triangle } a \ b \ z) \cup \text{path-inside } (\text{make-triangle } a \ b \ z)$

<proof>

lemma *smaller-triangle:*

assumes $\neg \text{collinear } \{a, b, c\} \wedge \neg \text{collinear } \{a', b', c'\}$

assumes $p = \text{make-triangle } a \ b \ c$

assumes $p' = \text{make-triangle } a' \ b' \ c'$

assumes $\text{path-inside } p \subseteq \text{path-inside } p'$

assumes $\exists d. \text{integral-vec } d \wedge d \in \text{path-image } p' \cup \text{path-inside } p' \wedge d \notin \text{path-image } p \cup \text{path-inside } p$

shows $\text{card } (\text{integral-inside } p) + \text{card } (\text{integral-boundary } p) < \text{card } (\text{integral-inside } p') + \text{card } (\text{integral-boundary } p')$

<proof>

lemma *pick-elem-triangle:*

fixes $p :: R\text{-to-}R^2$

assumes *p-triangle:* $p = \text{make-triangle } a \ b \ c$

assumes *elem-triangle:* $\text{elem-triangle } a \ b \ c$

assumes $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$ **and**

$B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$

shows $\text{measure lebesgue } (\text{path-inside } p) = I + B/2 - 1$

<proof>

lemma *pick-triangle-lemma:*

fixes $p :: R\text{-to-}R^2$

assumes $p = \text{make-triangle } a \ b \ c$ **and** *all-integral* $[a, b, c]$ **and** *distinct* $[a, b, c]$

and $\neg \text{collinear } \{a, b, c\}$

$I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$ **and**

$B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$

shows $\text{measure lebesgue } (\text{path-inside } p) = I + B/2 - 1$

<proof>

29.2 Pocket properties

definition *index-not-in-set* :: (real²) list ⇒ (real²) set ⇒ nat ⇒ bool
where *index-not-in-set vts A i* ↔ $i \in \{i. i < \text{length } vts \wedge vts ! i \notin A\}$

definition *min-index-not-in-set*:: (real²) list ⇒ (real²) set ⇒ nat
where *min-index-not-in-set vts A* = (LEAST *i. index-not-in-set vts A i*)

definition *nonzero-index-in-set* :: (real²) list ⇒ (real²) set ⇒ nat ⇒ bool
where
nonzero-index-in-set vts A i ↔ $i \in \{i. 0 < i \wedge i < \text{length } vts \wedge vts ! i \in A\}$

definition *min-nonzero-index-in-set* :: (real²) list ⇒ (real²) set ⇒ nat **where**
min-nonzero-index-in-set vts A = (LEAST *i. nonzero-index-in-set vts A i*)

definition *construct-pocket-0* :: (real²) list ⇒ (real²) set ⇒ (real²) list **where**
construct-pocket-0 vts A = take ((*min-nonzero-index-in-set vts A*) + 1) vts

definition *is-pocket-0* :: (real²) list ⇒ (real²) list ⇒ bool **where**
is-pocket-0 vts vts' ↔
 polygon (make-polygonal-path vts)
 ∧ (∃ *i. vts' = take i vts*)
 ∧ $3 \leq \text{length } vts' \wedge \text{length } vts' < \text{length } vts$
 ∧ *hd vts' ∈ frontier (convex hull (set vts))* ∧ *last vts' ∈ frontier (convex hull (set vts))*
 ∧ *set (tl (butlast vts')) ⊆ interior (convex hull (set vts))*

definition *fill-pocket-0* :: (real²) list ⇒ nat ⇒ (real²) list **where**
fill-pocket-0 vts i = (*hd vts*) # (*drop (i-1) vts*)

lemma *min-nonzero-index-in-set-exists*:
assumes *set (tl vts) ∩ A ≠ {}*
shows ∃ *i. nonzero-index-in-set vts A i*
 ⟨*proof*⟩

lemma *min-nonzero-index-in-set-defined*:
assumes *set (tl vts) ∩ A ≠ {}*
defines *i* ≡ *min-nonzero-index-in-set vts A*
shows *nonzero-index-in-set vts A i* ∧ (∀ *j < i. ¬ nonzero-index-in-set vts A j*)
 ⟨*proof*⟩

lemma *min-index-not-in-set-exists*:
assumes *set vts ⊃ A*
shows ∃ *i. index-not-in-set vts A i*
 ⟨*proof*⟩

lemma *min-index-not-in-set-defined*:
assumes *set vts ⊃ A*

defines $i \equiv \text{min-index-not-in-set } vts \ A$
shows $\text{index-not-in-set } vts \ A \ i \wedge (\forall j < i. \neg \text{index-not-in-set } vts \ A \ j)$
 $\langle \text{proof} \rangle$

lemma *min-nonzero-index-in-set-bound*:
assumes $\text{set } (tl \ vts) \cap A \neq \{\}$
shows $\text{min-nonzero-index-in-set } vts \ A < \text{length } vts$
 $\langle \text{proof} \rangle$

lemma *construct-pocket-0-subset-vts*:
assumes $\text{set } (tl \ vts) \cap A \neq \{\}$
shows $\text{set } (\text{construct-pocket-0 } vts \ A) \subseteq \text{set } vts$
 $\langle \text{proof} \rangle$

lemma *min-index-not-in-set-0*:
assumes $\text{set } vts \supset A$
assumes $vts!0 \in A$
defines $i \equiv \text{min-index-not-in-set } vts \ A$
defines $r \equiv i - 1$
shows $vts!r \in A$
 $\langle \text{proof} \rangle$

lemma *construct-pocket-0-last-in-set*:
assumes $\text{set } (tl \ vts) \cap A \neq \{\}$
assumes $vts!0 \in A$
defines $p \equiv \text{construct-pocket-0 } vts \ A$
shows $\text{last } p \in A$
 $\langle \text{proof} \rangle$

lemma *construct-pocket-0-first-last-distinct*:
assumes $\text{card } A \geq 2$
assumes $A \subseteq \text{set } vts$
assumes $\text{distinct } (\text{butlast } vts)$
assumes $\text{hd } vts = \text{last } vts$
shows $\text{hd } (\text{construct-pocket-0 } vts \ A) \neq \text{last } (\text{construct-pocket-0 } vts \ A)$
 $\langle \text{proof} \rangle$

lemma *construct-pocket-is-pocket*:
assumes $\text{polygon } (\text{make-polygonal-path } vts)$
assumes $vts!0 \in \text{frontier } (\text{convex hull } (\text{set } vts))$
assumes $vts!1 \notin \text{frontier } (\text{convex hull } (\text{set } vts))$
shows $\text{is-pocket-0 } vts \ (\text{construct-pocket-0 } vts \ (\text{set } vts \cap \text{frontier } (\text{convex hull } (\text{set } vts))))$
 $\langle \text{proof} \rangle$

lemma *exists-point-above-interior*:
fixes $a :: \text{real}^2$
assumes $a \in \text{interior } (\text{convex hull } S)$

obtains x **where** $x \in S \wedge x\$2 > a\2
 <proof>

lemma *exists-point-above-convex-hull-interior*:

fixes $S :: (\text{real}^2)$ *set*
assumes $S \neq \{\}$
assumes *compact* S
obtains x **where** $x \in S - (\text{interior } (\text{convex hull } S)) \wedge (\forall y \in \text{interior } (\text{convex hull } S). x\$2 > y\$2)$
 <proof>

lemma *flip-function*:

defines $M \equiv (\text{vector } [\text{vector } [1, 0], \text{vector } [0, -1]]) :: (\text{real}^2 \times \text{real}^2)$
defines $f \equiv \lambda v. M * v$
defines $g \equiv (\lambda v. \text{vector } [v\$1, -v\$2]) :: (\text{real}^2 \Rightarrow \text{real}^2)$
shows *inj* $f f = g$
 <proof>

lemma *exists-point-below-convex-hull-interior*:

fixes $S :: (\text{real}^2)$ *set*
assumes $S \neq \{\}$
assumes *compact* S
obtains x **where** $x \in S - (\text{interior } (\text{convex hull } S)) \wedge (\forall y \in \text{interior } (\text{convex hull } S). x\$2 < y\$2)$
 <proof>

lemma *exists-point-above-all*:

fixes $p\ q :: R\text{-to-}R^2$
defines $H \equiv \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
assumes *path* $p \wedge \text{path } q$
assumes $p\{0 < .. < 1\} \subseteq \text{interior } H$
assumes $(p\ 0)\$2 = 0 \wedge (p\ 1)\$2 = 0$
assumes $\exists x \in p\{0 < .. < 1\}. x\$2 \geq 0$
obtains x **where** $x \in \text{path-image } q \wedge (\forall y \in \text{path-image } p. x\$2 > y\$2)$
 <proof>

lemma *exists-point-below-all*:

fixes $p\ q :: R\text{-to-}R^2$
defines $H \equiv \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
assumes *path* $p \wedge \text{path } q$
assumes $p\{0 < .. < 1\} \subseteq \text{interior } H$
assumes $(p\ 0)\$2 = 0 \wedge (p\ 1)\$2 = 0$
assumes $\exists x \in \text{path-image } p \cup \text{path-image } q. x\$2 < 0$
obtains x **where** $x \in \text{path-image } q \wedge (\forall y \in \text{path-image } p. x\$2 < y\$2)$
 <proof>

lemma *pocket-fill-line-int-aux*:

fixes $x\ y\ z :: \text{real}^2$
defines $a \equiv y\$1$

assumes $x = 0$
assumes $a > 0 \wedge y\$2 = 0$
assumes $z\$1 < 0 \vee z\$1 > a$
assumes $z\$2 = 0$
assumes $\text{convex } A \wedge \text{compact } A$
assumes $\{x, y, z\} \subseteq A$
assumes $\{x, y\} \subseteq \text{frontier } A$
shows $z \in \text{frontier } A \wedge \text{closed-segment } x y \subseteq \text{frontier } A$
 <proof>

lemma *axis-dist*:

fixes $a b :: \text{real}^2$
shows $a\$2 = b\$2 \implies \text{dist } a b = \text{dist } (a\$1) (b\$1) \quad a\$1 = b\$1 \implies \text{dist } a b = \text{dist } (a\$2) (b\$2)$
 <proof>

lemma *dist-bound-1*:

fixes $a b x :: \text{real}^2$
assumes $a\$2 = x\2
assumes $b \in \text{ball } x \ \varepsilon$
assumes $\varepsilon < \text{dist } a x$
shows $a\$1 < x\$1 \implies b\$1 > a\$1 \quad a\$1 > x\$1 \implies b\$1 < a\1
 <proof>

lemma *dist-bound-2*:

fixes $a b x :: \text{real}^2$
assumes $a\$1 = x\1
assumes $b \in \text{ball } x \ \varepsilon$
assumes $\varepsilon < \text{dist } a x$
shows $a\$2 < x\$2 \implies b\$2 > a\$2 \quad a\$2 > x\$2 \implies b\$2 < a\2
 <proof>

lemma *linepath-bound-1*:

fixes $x y :: \text{real}^2$
shows $a < x\$1 \wedge a < y\$1 \implies \forall v \in \text{path-image } (\text{linepath } x y). a < v\1
 $x\$1 < b \wedge y\$1 < b \implies \forall v \in \text{path-image } (\text{linepath } x y). v\$1 < b$
 <proof>

lemma *linepath-bound-2*:

fixes $x y :: \text{real}^2$
shows $a < x\$2 \wedge a < y\$2 \implies \forall v \in \text{path-image } (\text{linepath } x y). a < v\2
 $x\$2 < b \wedge y\$2 < b \implies \forall v \in \text{path-image } (\text{linepath } x y). v\$2 < b$
 <proof>

lemma *linepath-int-corner*:

fixes $x y z :: \text{real}^2$
assumes $x\$2 \neq y\2
assumes $y\$2 = z\2
shows $\text{path-image } (\text{linepath } x y) \cap \text{path-image } (\text{linepath } y z) = \{y\}$

(is path-image ?l1 \cap path-image ?l2 = {y})
 <proof>

lemma *linepath-int-vertical*:

fixes $w\ x\ y\ z :: \text{real}^2$
assumes $w\$1 \neq y\1
assumes $w\$1 = x\1
assumes $y\$1 = z\1
shows $\text{path-image } (\text{linepath } w\ x) \cap \text{path-image } (\text{linepath } y\ z) = \{\}$
 <proof>

lemma *linepath-int-horizontal*:

fixes $w\ x\ y\ z :: \text{real}^2$
assumes $w\$2 \neq y\2
assumes $w\$2 = x\2
assumes $y\$2 = z\2
shows $\text{path-image } (\text{linepath } w\ x) \cap \text{path-image } (\text{linepath } y\ z) = \{\}$
 <proof>

lemma *linepath-int-columns*:

fixes $w\ x\ y\ z :: \text{real}^2$
assumes $w\$1 < y\$1 \wedge w\$1 < z\1
assumes $x\$1 < y\$1 \wedge x\$1 < z\1
shows $\text{path-image } (\text{linepath } w\ x) \cap \text{path-image } (\text{linepath } y\ z) = \{\}$
 (is path-image ?l1 \cap path-image ?l2 = {})
 <proof>

lemma *linepath-int-rows*:

fixes $w\ x\ y\ z :: \text{real}^2$
assumes $w\$2 < y\$2 \wedge w\$2 < z\2
assumes $x\$2 < y\$2 \wedge x\$2 < z\2
shows $\text{path-image } (\text{linepath } w\ x) \cap \text{path-image } (\text{linepath } y\ z) = \{\}$
 (is path-image ?l1 \cap path-image ?l2 = {})
 <proof>

lemma *horizontal-segment-at-0*:

assumes $a > 0$
shows $\text{closed-segment } ((\text{vector } [0, 0])::(\text{real}^2)) (\text{vector } [a, 0]) = \{x. x\$2 = 0 \wedge x\$1 \in \{0..a\}\}$
 (is ?l = ?s)
 <proof>

lemma *horizontal-segment-at-0'*:

fixes $x\ y :: \text{real}^2$
assumes $a > 0$
assumes $x\$1 = 0 \wedge x\$2 = 0 \wedge y\$1 = a \wedge y\$2 = 0$
shows $\text{closed-segment } x\ y = \{x. x\$2 = 0 \wedge x\$1 \in \{0..a\}\}$
 <proof>

lemma *pocket-fill-line-int-aux1*:
fixes $p\ q :: R\text{-to-}R^2$
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $q0 \equiv \text{pathstart } q$
defines $q1 \equiv \text{pathfinish } q$
defines $a \equiv p1\$1$
defines $l \equiv \text{closed-segment } p0\ p1$
assumes *simple-path* p
assumes *simple-path* q
assumes $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$
assumes $a > 0$
assumes $\text{path-image } q \cap \{x. x\$2 = 0\} \subseteq l$
assumes $\text{path-image } p \cap \{x. x\$2 = 0\} \subseteq l$
assumes $\forall v \in \text{path-image } p. q0\$2 \leq v\$2$
assumes $\forall v \in \text{path-image } p. q1\$2 > v\$2$
shows $\text{path-image } p \cap \text{path-image } q \neq \{\}$
<proof>

lemma *pocket-fill-line-int-aux2*:
fixes $p\ q :: R\text{-to-}R^2$
fixes $A :: (\text{real}^2)\ \text{set}$
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $a \equiv p1\$1$
defines $l \equiv \text{closed-segment } p0\ p1$
assumes *simple-path* p
assumes $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$
assumes $a > 0$
assumes *convex* $A \wedge \text{compact } A$
assumes $\{p0, p1\} \subseteq \text{frontier } A$
assumes $p \text{ ' } \{0 <..<1\} \subseteq \text{interior } A$
shows $\text{path-image } p \cap \{x. x\$2 = 0\} \subseteq l$
<proof>

lemma *three-points-on-line*:
fixes $a\ b :: 'a::\text{real-vector}$
assumes $A = \text{affine hull } \{a, b\}$
assumes $a \neq b$
assumes $\{x, y, z\} \subseteq A$
assumes $x \neq y \wedge y \neq z \wedge x \neq z$
shows $x \in \text{open-segment } y\ z \vee y \in \text{open-segment } x\ z \vee z \in \text{open-segment } x\ y$
<proof>

lemma *pocket-fill-line-int-aux3*:
fixes $A :: (\text{real}^2)\ \text{set}$
assumes *convex* $A \wedge \text{compact } A$
assumes $v \neq 0$
assumes $\text{closed-segment } 0\ w \subseteq \text{frontier } A$ (**is** *closed-segment* $?a\ ?b \subseteq -$)

assumes $w \cdot v = 0$
assumes $w \neq 0$
shows $(A \subseteq \{x. x \cdot v \leq 0\} \vee A \subseteq \{x. x \cdot v \geq 0\})$ (**is** $A \subseteq ?P1 \vee A \subseteq ?P2$)
 <proof>

lemma *pocket-fill-line-int-aux4*:

fixes $p\ q :: R\text{-to-}R^2$
fixes $A :: (\text{real}^2)$ set
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $q0 \equiv \text{pathstart } q$
defines $q1 \equiv \text{pathfinish } q$
defines $a \equiv p1\$1$
defines $l \equiv \text{closed-segment } p0\ p1$
assumes *simple-path* p
assumes *simple-path* q
assumes $\text{path-image } p \cap \text{path-image } q = \{\}$
assumes $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$
assumes $a > 0$
assumes $\forall v \in \text{path-image } p. q0\$2 \leq v\$2$
assumes $\forall v \in \text{path-image } p. q1\$2 > v\$2$
assumes *convex* $A \wedge$ *compact* A
assumes $\{p0, p1\} \subseteq \text{frontier } A$
assumes $p\{0<..
assumes $\text{path-image } q \subseteq A$
shows $l \subseteq \text{frontier } A \forall x \in (\text{path-image } p) \cup (\text{path-image } q). x\$2 \geq 0 \wedge q0\$2 = 0$
 <proof>$

lemma *pocket-fill-line-int-aux5*:

fixes $p\ q :: R\text{-to-}R^2$
fixes $A :: (\text{real}^2)$ set
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $q0 \equiv \text{pathstart } q$
defines $q1 \equiv \text{pathfinish } q$
defines $a \equiv p1\$1$
defines $l \equiv \text{closed-segment } p0\ p1$
assumes *simple-path* p
assumes *simple-path* q
assumes $\text{path-image } p \cap \text{path-image } q = \{q0, q1\}$
assumes $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$
assumes $a > 0$
assumes $A = \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
assumes $\{p0, p1\} \subseteq \text{frontier } A$
assumes $p\{0<..
assumes $\text{path-image } q \subseteq A$
assumes $\exists x \in p\{0<..
assumes $q0 = p1 \wedge q1 = p0$$$

shows $l \subseteq \text{frontier } A \ \forall x \in \text{path-image } p \cup \text{path-image } q. \ x \geq 0$
 <proof>

lemma *pocket-fill-line-int-aux6*:

fixes $p \ q :: R\text{-to-}R^2$
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $q0 \equiv \text{pathstart } q$
defines $q1 \equiv \text{pathfinish } q$
defines $a \equiv p1 \cdot 1$
assumes *simple-path* p
assumes *simple-path* q
assumes $p0 = 0 \wedge p1 \cdot 2 = 0$
assumes $a > 0$
assumes $q0 \cdot 1 \in \{0..a\} \wedge q0 \cdot 2 = 0$
assumes $\forall x \in \text{path-image } p. \ q1 \cdot 2 > x \cdot 2$
assumes $\forall x \in \text{path-image } p \cup \text{path-image } q. \ x \geq 0$
shows $\text{path-image } p \cap \text{path-image } q \neq \{\}$
 <proof>

lemma *pocket-fill-line-int-aux7*:

fixes $p \ q :: R\text{-to-}R^2$
fixes $A :: (\text{real}^2) \text{ set}$
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $q0 \equiv \text{pathstart } q$
defines $q1 \equiv \text{pathfinish } q$
defines $a \equiv p1 \cdot 1$
defines $l \equiv \text{open-segment } p0 \ p1$
assumes *simple-path* p
assumes *simple-path* q
assumes $\text{path-image } p \cap \text{path-image } q = \{q0, q1\}$
assumes $p0 \cdot 1 = 0 \wedge p0 \cdot 2 = 0 \wedge p1 \cdot 2 = 0$
assumes $a > 0$
assumes $A = \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
assumes $\{p0, p1\} \subseteq \text{frontier } A$
assumes $p \cdot \{0 < .. < 1\} \subseteq \text{interior } A$
assumes $\exists x \in p \cdot \{0 < .. < 1\}. \ x \geq 0$
assumes $q0 = p1 \wedge q1 = p0$
shows $\text{path-image } q \cap l = \{\}$ *closed-segment* $p0 \ p1 \subseteq \text{frontier } A$
 <proof>

lemma *frontier-injective-linear-image*:

fixes $f :: 'a::\text{euclidean-space} \Rightarrow 'a::\text{euclidean-space}$
assumes *linear* f *inj* f
shows $f \cdot (\text{frontier } S) = \text{frontier } (f \cdot S)$
 <proof>

lemma *pocket-fill-line-int-aux8*:
fixes $p\ q :: R\text{-to-}R^2$
fixes $A :: (\text{real}^2)\ \text{set}$
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $q0 \equiv \text{pathstart } q$
defines $q1 \equiv \text{pathfinish } q$
defines $a \equiv p1\$1$
defines $l \equiv \text{open-segment } p0\ p1$
assumes *simple-path* p
assumes *simple-path* q
assumes $\text{path-image } p \cap \text{path-image } q = \{q0, q1\}$
assumes $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$
assumes $a > 0$
assumes $A = \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
assumes $\{p0, p1\} \subseteq \text{frontier } A$
assumes $p\{0 < .. < 1\} \subseteq \text{interior } A$
assumes $q0 = p1 \wedge q1 = p0$
shows $\text{path-image } q \cap l = \{\} \wedge l \subseteq \text{frontier } A$
 $\langle \text{proof} \rangle$

lemma *simple-path-linear-image*:
assumes *simple-path* p
assumes *inj* $f \wedge \text{bounded-linear } f$
shows *simple-path* $(f \circ p)$
 $\langle \text{proof} \rangle$

lemma *pts-interior*:
fixes pts
defines $p \equiv \text{make-polygonal-path } pts$
assumes *convex* H
assumes $\forall j \in \{0 < .. < \text{length } pts - 1\}. pts!j \notin \text{frontier } H$
assumes *loop-free* p
assumes $\text{path-image } p \subseteq H$
assumes $\text{length } pts \geq 3$
shows $p\{0 < .. < 1\} \subseteq \text{interior } H$
 $\langle \text{proof} \rangle$

lemma *pocket-fill-line-int-0*:
assumes *polygon-of* $r\ pts$
defines $H \equiv \text{convex hull } (\text{set } pts)$
assumes $2 \leq i \wedge i < \text{length } pts - 1$
defines $a \equiv \text{hd } pts$
defines $b \equiv pts!i$
assumes $\{a, b\} \subseteq \text{frontier } H$
assumes $\forall j \in \{0 < .. < i\}. pts!j \notin \text{frontier } H$
assumes $a = 0$
shows $\text{path-image } (\text{linepath } a\ b) \cap \text{path-image } r = \{a, b\}$
 $\text{path-image } (\text{linepath } a\ b) \subseteq \text{frontier } H$

<proof>

lemma *linepath-translation*: $(\lambda v. v - a) \circ (\text{linepath } x \ y) = \text{linepath } ((\lambda v. v - a) \ x) \ ((\lambda v. v - a) \ y)$
<proof>

lemma *linepath-image-translation*:
 $\text{path-image } ((\lambda v. v - a) \circ (\text{linepath } x \ y)) = \text{path-image } (\text{linepath } ((\lambda v. v - a) \ x) \ ((\lambda v. v - a) \ y))$
<proof>

lemma *make-polygonal-path-translate*:
assumes $\text{length } vts \geq 1$
shows $(\lambda v. v - a) \circ (\text{make-polygonal-path } vts) = \text{make-polygonal-path } (\text{map } (\lambda v. v - a) \ vts)$
<proof>

lemma *pocket-fill-line-int*:
assumes *polygon-of* $r \ vts$
defines $H \equiv \text{convex hull } (\text{set } vts)$
assumes $2 \leq i \wedge i < \text{length } vts - 1$
defines $a \equiv \text{hd } vts$
defines $b \equiv vts!i$
assumes $\{a, b\} \subseteq \text{frontier } H$
assumes $\forall j \in \{0 <..<i\}. vts!j \notin \text{frontier } H$
shows $\text{path-image } (\text{linepath } a \ b) \cap \text{path-image } r = \{a, b\}$
 $\text{path-image } (\text{linepath } a \ b) \subseteq \text{frontier } H$
<proof>

lemma *path-connected-simple-path-endless*:
assumes *simple-path* p
shows *path-connected* $(\text{path-image } p - \{\text{pathstart } p, \text{pathfinish } p\})$ (**is** *path-connected* $?S$)
<proof>

lemma *simple-loop-split*:
assumes *simple-path* $p \wedge \text{closed-path } p$
assumes *simple-path* q
assumes $\text{path-image } q \cap \text{path-image } p = \{q \ 0, q \ 1\}$
assumes $\text{path-image } q \cap \text{path-inside } p \neq \{\}$
shows $q\{0 <..<1\} \subseteq \text{path-inside } p$
<proof>

lemma *pocket-path-interior-arc*:
assumes *simple-path* $p \wedge \text{simple-path } q$
assumes *arc* $p \wedge \text{arc } q$
assumes $q \ 0 = p \ 1 \wedge q \ 1 = p \ 0$
assumes $\text{path-image } p \cap \text{path-image } q = \{p \ 0, q \ 0\}$

defines $A \equiv \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
defines $l \equiv \text{linepath } (p\ 0) (p\ 1)$
assumes $p\{0 < .. < 1\} \subseteq \text{interior } A$
assumes $\text{path-image } l \subseteq \text{frontier } A$
assumes $\text{path-image } q \cap \text{path-image } l = \{l\ 0, q\ 0\}$
shows $p\{0 < .. < 1\} \cap \text{path-inside } (l\ +++\ q) \neq \{\}$
 $\text{simple-path } (l\ +++\ q) \wedge \text{closed-path } (l\ +++\ q)$
 $\text{path-image } p \cap \text{path-image } (l\ +++\ q) = \{p\ 0, p\ 1\}$
<proof>

lemma *pocket-path-interior*:

assumes $\text{simple-path } p \wedge \text{simple-path } q$
assumes $\text{arc } p \wedge \text{arc } q$
assumes $q\ 0 = p\ 1 \wedge q\ 1 = p\ 0$
assumes $\text{path-image } p \cap \text{path-image } q = \{p\ 0, q\ 0\}$
defines $A \equiv \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
defines $l \equiv \text{linepath } (p\ 0) (p\ 1)$
assumes $p\{0 < .. < 1\} \subseteq \text{interior } A$
assumes $\text{path-image } l \subseteq \text{frontier } A$
assumes $\text{path-image } q \cap \text{path-image } l = \{l\ 0, q\ 0\}$
shows $p\{0 < .. < 1\} \subseteq \text{path-inside } (l\ +++\ q)$
<proof>

lemma *pocket-path-good*:

assumes $\text{polygon } (\text{make-polygonal-path } vts)$
assumes $vts!0 \in \text{frontier } (\text{convex hull } (\text{set } vts))$
assumes $vts!1 \notin \text{frontier } (\text{convex hull } (\text{set } vts))$
assumes $\neg \text{convex } (\text{path-image } (\text{make-polygonal-path } vts) \cup \text{path-inside } (\text{make-polygonal-path } vts))$
defines $\text{pocket-path-vts} \equiv \text{construct-pocket-0 } vts (\text{set } vts \cap \text{frontier } (\text{convex hull } (\text{set } vts)))$
defines $\text{pocket} \equiv \text{make-polygonal-path } (\text{pocket-path-vts} @ [\text{pocket-path-vts!0}])$
defines $\text{filled-vts} \equiv \text{fill-pocket-0 } vts (\text{length } \text{pocket-path-vts})$
defines $\text{filled-p} \equiv \text{make-polygonal-path } \text{filled-vts}$
defines $a \equiv \text{hd } \text{pocket-path-vts}$
defines $b \equiv \text{last } \text{pocket-path-vts}$
defines $\text{good-pocket-path-vts} \equiv \text{tl } (\text{butlast } \text{pocket-path-vts})$
shows $\text{polygon } \text{filled-p}$
 $\text{is-polygon-split-path } (\text{butlast } \text{filled-vts})\ 0\ 1\ \text{good-pocket-path-vts}$
 $\text{polygon } \text{pocket}$
 $\text{card } (\text{set } \text{pocket-path-vts}) < \text{card } (\text{set } vts)$
 $\text{card } (\text{set } \text{filled-vts}) < \text{card } (\text{set } vts)$
<proof>

29.3 Arbitrary Polygon Case

lemma *pick-rotate*:

assumes $\text{polygon-of } p\ vts$
assumes $\text{all-integral } vts$

obtains $p' \text{ vts}'$ **where** *polygon-of* $p' \text{ vts}'$
 $\wedge \text{vts}'!0 \in \text{frontier} (\text{convex hull} (\text{set } \text{vts}'))$
 $\wedge \text{path-image } p' = \text{path-image } p$
 $\wedge \text{all-integral } \text{vts}'$
 $\wedge \text{set } \text{vts}' = \text{set } \text{vts}$
 <proof>

lemma *pick-unrotated*:
fixes $p :: R\text{-to-}R^2$
assumes *polygon*: *polygon* p
assumes *polygonal-path*: $p = \text{make-polygonal-path } \text{vts}$
assumes *int-vertices*: *all-integral* vts
assumes *I-is*: $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$
assumes *B-is*: $B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$
assumes $\text{vts}'!0 \in \text{frontier} (\text{convex hull} (\text{set } \text{vts}'))$
shows *measure lebesgue* (*path-inside* p) = $I + B/2 - 1$
 <proof>

theorem *pick*:
fixes $p :: R\text{-to-}R^2$
assumes *polygon* p
assumes $p = \text{make-polygonal-path } \text{vts}$
assumes *all-integral* vts
assumes $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$
assumes $B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$
shows *measure lebesgue* (*path-inside* p) = $I + B/2 - 1$
 <proof>

end

References

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