

Perron-Frobenius Theorem for Spectral Radius Analysis*

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Abstract

The spectral radius of a matrix A is the maximum norm of all eigenvalues of A . In previous work we already formalized that for a complex matrix A , the values in A^n grow polynomially in n if and only if the spectral radius is at most one. One problem with the above characterization is the determination of all *complex* eigenvalues. In case A contains only non-negative real values, a simplification is possible with the help of the Perron-Frobenius theorem, which tells us that it suffices to consider only the *real* eigenvalues of A , i.e., applying Sturm's method can decide the polynomial growth of A^n .

We formalize the Perron-Frobenius theorem based on a proof via Brouwer's fixpoint theorem, which is available in the HOL multivariate analysis (HMA) library. Since the results on the spectral radius is based on matrices in the Jordan normal form (JNF) library, we further develop a connection which allows us to easily transfer theorems between HMA and JNF. With this connection we derive the combined result: if A is a non-negative real matrix, and no real eigenvalue of A is strictly larger than one, then A^n is polynomially bounded in n .

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1 Introduction

The spectral radius of a matrix A over \mathbb{R} or \mathbb{C} is defined as

$$\rho(A) = \max \{|x| \mid \chi_A(x) = 0, x \in \mathbb{C}\}$$

where χ_A is the characteristic polynomial of A . It is a central notion related to the growth rate of matrix powers. A matrix A has polynomial growth, i.e., all values of A^n can be bounded polynomially in n , if and only if $\rho(A) \leq 1$. It is quite easy to see that $\rho(A) \leq 1$ is a necessary criterion,¹ but it is more complicated to argue about sufficiency. In previous work we formalized this statement via Jordan normal forms [4].

Theorem 1 (in JNF). *The values in A^n are polynomially bounded in n if $\rho(A) \leq 1$.*

In order to perform the proof via Jordan normal forms, we did not use the HMA library from the distribution to represent matrices. The reason is that already the definition of a Jordan normal form is naturally expressed via block-matrices, and arbitrary block-matrices are hard to express in HMA, if at all.

¹Let λ and v be some eigenvalue and eigenvector pair such that $|\lambda| > 1$. Then $|A^n v| = |\lambda^n v| = |\lambda|^n |v|$ grows exponentially in n , where $|w|$ denotes the component-wise application of $|\cdot|$ to vector elements of w .

The problem in applying Theorem 1 in concrete examples is the determination of all complex roots of the polynomial χ_A . For instance, one can utilize complex algebraic numbers for this purpose, which however are computationally expensive. To avoid this problem, in this work we formalize the Perron Frobenius theorem. It states that for non-negative real-valued matrices, $\rho(A)$ is an eigenvalue of A .

Theorem 2 (in HMA). *If $A \in \mathbb{R}_{\geq 0}^{k \times k}$, then $\chi_A(\rho(A)) = 0$.*

We decided to perform the formalization based on the HMA library, since there is a short proof of Theorem 2 via Brouwer’s fixpoint theorem [2, Section 5.2]. The latter is a well-known but complex theorem that is available in HMA, but not in the JNF library.

Eventually we want to combine both theorems to obtain:

Corollary 1. *If $A \in \mathbb{R}_{\geq 0}^{k \times k}$, then the values in A^n are polynomially bounded in n if χ_A has no real roots in the interval $(1, \infty)$.*

This criterion is computationally far less expensive – one invocation of Sturm’s method on χ_A suffices. Unfortunately, we cannot immediately combine both theorems. We first have to bridge the gap between the HMA-world and the JNF-world. To this end, we develop a setup for the transfer-tool which admits to translate theorems from JNF into HMA. Moreover, using a recent extension for local type definitions within proofs [1], we also provide a translation from HMA into JNF.

With the help of these translations, we prove Corollary 1 and make it available in both HMA and JNF. (In the formalization the corollary looks a bit more complicated as it also contains an estimation of the the degree of the polynomial growth.)

2 Elimination of CARD('n)

In the following theory we provide a method which modifies theorems of the form $P[\text{CARD}(n)]$ into $n! = 0 \implies P[n]$, so that they can more easily be applied.

Known issues: there might be problems with nested meta-implications and meta-quantification.

theory *Cancel-Card-Constraint*

imports

HOL-Types-To-Sets.Types-To-Sets

HOL-Library.Cardinality

begin

lemma *n-zero-nonempty*: $n \neq 0 \implies \{0 ..< n :: nat\} \neq \{\}$ **by** *auto*

```

lemma type-impl-card-n: assumes  $\exists (Rep :: 'a \Rightarrow nat) Abs. \text{type-definition } Rep$ 
Abs  $\{0 ..< n :: nat\}$ 
  shows class.finite (TYPE('a))  $\wedge$  CARD('a) = n
proof –
  from assms obtain rep :: 'a  $\Rightarrow$  nat and abs :: nat  $\Rightarrow$  'a where t: type-definition
rep abs  $\{0 ..< n\}$  by auto
  have card (UNIV :: 'a set) = card  $\{0 ..< n\}$  using t by (rule type-definition.card)
  also have ... = n by auto
  finally have bn: CARD ('a) = n .
  have finite (abs '  $\{0 ..< n\}$ ) by auto
  also have abs '  $\{0 ..< n\}$  = UNIV using t by (rule type-definition.Abs-image)
  finally have class.finite (TYPE('a)) unfolding class.finite-def .
  with bn show ?thesis by blast
qed

```

ML-file \langle *cancel-card-constraint.ML* \rangle

end

3 Connecting HMA-matrices with JNF-matrices

The following theories provide a connection between the type-based representation of vectors and matrices in HOL multivariate-analysis (HMA) with the set-based representation of vectors and matrices with integer indices in the Jordan-normal-form (JNF) development.

3.1 Bijections between index types of HMA and natural numbers

At the core of HMA-connect, there has to be a translation between indices of vectors and matrices, which are via index-types on the one hand, and natural numbers on the other hand.

We some unspecified bijection in our application, and not the conversions to-nat and from-nat in theory Rank-Nullity-Theorem/Mod-Type, since our definitions below do not enforce any further type constraints.

```

theory Bij-Nat
imports
  HOL-Library.Cardinality
  HOL-Library.Natural-Number
begin

```

```

lemma finite-set-to-list:  $\exists xs :: 'a :: \text{finite list. } \text{distinct } xs \wedge \text{set } xs = Y$ 

```

```

proof –
  have finite Y by simp
  thus ?thesis
  proof (induct Y rule: finite-induct)
    case (insert y Y)
    then obtain xs where xs: distinct xs set xs = Y by auto
    show ?case
      by (rule exI[of - y # xs], insert xs insert(2), auto)
  qed simp
qed

```

```

definition univ-list :: 'a :: finite list where
  univ-list = (SOME xs. distinct xs ∧ set xs = UNIV)

```

```

lemma univ-list: distinct (univ-list :: 'a list) set univ-list = (UNIV :: 'a :: finite set)

```

```

proof –
  let ?xs = univ-list :: 'a list
  have distinct ?xs ∧ set ?xs = UNIV
    unfolding univ-list-def
    by (rule someI-ex, rule finite-set-to-list)
  thus distinct ?xs set ?xs = UNIV by auto
qed

```

```

definition to-nat :: 'a :: finite ⇒ nat where
  to-nat a = (SOME i. univ-list ! i = a ∧ i < length (univ-list :: 'a list))

```

```

definition from-nat :: nat ⇒ 'a :: finite where
  from-nat i = univ-list ! i

```

```

lemma length-univ-list-card: length (univ-list :: 'a :: finite list) = CARD('a)
  using distinct-card[of univ-list :: 'a list, symmetric]
  by (auto simp: univ-list)

```

```

lemma to-nat-ex: ∃! i. univ-list ! i = (a :: 'a :: finite) ∧ i < length (univ-list :: 'a list)

```

```

proof –
  let ?ul = univ-list :: 'a list
  have a-in-set: a ∈ set ?ul unfolding univ-list by auto
  from this [unfolded set-conv-nth]
  obtain i where i1: ?ul ! i = a ∧ i < length ?ul by auto
  show ?thesis
  proof (rule ex1I, rule i1)
    fix j
    assume ?ul ! j = a ∧ j < length ?ul
    moreover have distinct ?ul by (simp add: univ-list)
    ultimately show j = i using i1 nth-eq-iff-index-eq by blast
  qed
qed

```

lemma *to-nat-less-card*: $to\text{-}nat\ (a :: 'a :: finite) < CARD('a)$
proof –
let $?ul = univ\text{-}list :: 'a\ list$
from *to-nat-ex*[*of a*] **obtain** i **where**
 $i1: univ\text{-}list ! i = a \wedge i < length\ (univ\text{-}list :: 'a\ list)$ **by** *auto*
show *?thesis* **unfolding** *to-nat-def*
proof (*rule someI2, rule i1*)
fix x
assume $x: ?ul ! x = a \wedge x < length\ ?ul$
thus $x < CARD\ ('a)$ **using** x **by** (*simp add: univ-list length-univ-list-card*)
qed
qed

lemma *to-nat-from-nat-id*:
assumes $i: i < CARD('a :: finite)$
shows $to\text{-}nat\ (from\text{-}nat\ i :: 'a) = i$
unfolding *to-nat-def from-nat-def*
proof (*rule some-equality, simp*)
have $l: length\ (univ\text{-}list :: 'a\ list) = card\ (set\ (univ\text{-}list :: 'a\ list))$
by (*rule distinct-card[symmetric], simp add: univ-list*)
thus $i2: i < length\ (univ\text{-}list :: 'a\ list)$
using i **unfolding** *univ-list* **by** *simp*
fix n
assume $n: (univ\text{-}list :: 'a\ list) ! n = (univ\text{-}list :: 'a\ list) ! i \wedge n < length\ (univ\text{-}list :: 'a\ list)$
have $d: distinct\ (univ\text{-}list :: 'a\ list)$ **using** *univ-list* **by** *simp*
show $n = i$ **using** *nth-eq-iff-index-eq[OF d - i2]* n **by** *auto*
qed

lemma *from-nat-inj*: **assumes** $i: i < CARD('a :: finite)$
and $j: j < CARD('a :: finite)$
and $id: (from\text{-}nat\ i :: 'a) = from\text{-}nat\ j$
shows $i = j$
proof –
from *arg-cong[OF id, of to-nat]*
show *?thesis* **using** $i\ j$ **by** (*simp add: to-nat-from-nat-id*)
qed

lemma *from-nat-to-nat-id[simp]*:
 $(from\text{-}nat\ (to\text{-}nat\ a)) = (a :: 'a :: finite)$
proof –
have $a\text{-in-set}: a \in set\ (univ\text{-}list)$ **unfolding** *univ-list* **by** *auto*
from *this* [*unfolded set-conv-nth*]
obtain i **where** $i1: univ\text{-}list ! i = a \wedge i < length\ (univ\text{-}list :: 'a\ list)$ **by** *auto*
show *?thesis*
unfolding *to-nat-def from-nat-def*
by (*rule someI2, rule i1, simp*)
qed

```

lemma to-nat-inj[simp]: assumes to-nat a = to-nat b
  shows a = b
proof -
  from to-nat-ex[of a] to-nat-ex[of b]
  show a = b unfolding to-nat-def by (metis assms from-nat-to-nat-id)
qed

lemma range-to-nat: range (to-nat :: 'a :: finite  $\Rightarrow$  nat) = {0 ..< CARD('a)} (is
?l = ?r)
proof -
  {
    fix i
    assume i  $\in$  ?l
    hence i  $\in$  ?r using to-nat-less-card[where 'a = 'a] by auto
  }
  moreover
  {
    fix i
    assume i  $\in$  ?r
    hence i < CARD('a) by auto
    from to-nat-from-nat-id[OF this]
    have i  $\in$  ?l by (metis range-eqI)
  }
  ultimately show ?thesis by auto
qed

lemma inj-to-nat: inj to-nat by (simp add: inj-on-def)

lemma bij-to-nat: bij-betw to-nat (UNIV :: 'a :: finite set) {0 ..< CARD('a)}
  unfolding bij-betw-def by (auto simp: range-to-nat inj-to-nat)

lemma numeral-nat: (numeral m1 :: nat) * numeral n1  $\equiv$  numeral (m1 * n1)
  (numeral m1 :: nat) + numeral n1  $\equiv$  numeral (m1 + n1) by simp-all

lemmas card-num-simps =
  card-num1 card-bit0 card-bit1
  mult-num-simps
  add-num-simps
  eq-num-simps
  mult-Suc-right mult-0-right One-nat-def add.right-neutral
  numeral-nat Suc-numeral

end

```

3.2 Transfer rules to convert theorems from JNF to HMA and vice-versa.

theory *HMA-Connect*

imports

Jordan-Normal-Form.Spectral-Radius
HOL-Analysis.Determinants
HOL-Analysis.Cartesian-Euclidean-Space
Bij-Nat
Cancel-Card-Constraint
HOL-Eisbach.Eisbach

begin

Prefer certain constants and lemmas without prefix.

hide-const (**open**) *Matrix.mat*

hide-const (**open**) *Matrix.row*

hide-const (**open**) *Determinant.det*

lemmas *mat-def* = *Finite-Cartesian-Product.mat-def*

lemmas *det-def* = *Determinants.det-def*

lemmas *row-def* = *Finite-Cartesian-Product.row-def*

notation *vec-index* (**infixl** \$v 90\$)

notation *vec-nth* (**infixl** \$h 90\$)

Forget that *'a mat*, *'a Matrix.vec*, and *'a poly* have been defined via lifting

lifting-forget *vec.lifting*

lifting-forget *mat.lifting*

lifting-forget *poly.lifting*

Some notions which we did not find in the HMA-world.

definition *eigen-vector* :: *'a::comm-ring-1 ^ 'n ^ 'n* \Rightarrow *'a ^ 'n* \Rightarrow *'a* \Rightarrow *bool* **where**
eigen-vector *A v ev* = (*v* \neq 0 \wedge *A* * *v* = *ev* * *s* *v*)

definition *eigen-value* :: *'a :: comm-ring-1 ^ 'n ^ 'n* \Rightarrow *'a* \Rightarrow *bool* **where**
eigen-value *A k* = (\exists *v*. *eigen-vector* *A v k*)

definition *similar-matrix-wit*

:: *'a :: semiring-1 ^ 'n ^ 'n* \Rightarrow *'a ^ 'n ^ 'n* \Rightarrow *'a ^ 'n ^ 'n* \Rightarrow *'a ^ 'n ^ 'n* \Rightarrow *bool*

where

similar-matrix-wit *A B P Q* = (*P* ** *Q* = *mat 1* \wedge *Q* ** *P* = *mat 1* \wedge *A* = *P* ** *B* ** *Q*)

definition *similar-matrix*

:: *'a :: semiring-1 ^ 'n ^ 'n* \Rightarrow *'a ^ 'n ^ 'n* \Rightarrow *bool* **where**

similar-matrix *A B* = (\exists *P Q*. *similar-matrix-wit* *A B P Q*)

definition *spectral-radius* :: $\text{complex } ^\wedge 'n \ ^\wedge 'n \Rightarrow \text{real}$ **where**
spectral-radius $A = \text{Max } \{ \text{norm } ev \mid v \text{ ev. eigen-vector } A \ v \text{ ev} \}$

definition *Spectrum* :: $'a :: \text{field } ^\wedge 'n \ ^\wedge 'n \Rightarrow 'a \text{ set}$ **where**
Spectrum $A = \text{Collect } (\text{eigen-value } A)$

definition *vec-elements-h* :: $'a \ ^\wedge 'n \Rightarrow 'a \text{ set}$ **where**
vec-elements-h $v = \text{range } (\text{vec-nth } v)$

lemma *vec-elements-h-def'*: $\text{vec-elements-h } v = \{v \ \$h \ i \mid i. \text{True}\}$
unfolding *vec-elements-h-def* **by** *auto*

definition *elements-mat-h* :: $'a \ ^\wedge 'nc \ ^\wedge 'nr \Rightarrow 'a \text{ set}$ **where**
elements-mat-h $A = \text{range } (\lambda \ (i,j). \ A \ \$h \ i \ \$h \ j)$

lemma *elements-mat-h-def'*: $\text{elements-mat-h } A = \{A \ \$h \ i \ \$h \ j \mid i \ j. \ \text{True}\}$
unfolding *elements-mat-h-def* **by** *auto*

definition *map-vector* :: $('a \Rightarrow 'b) \Rightarrow 'a \ ^\wedge 'n \Rightarrow 'b \ ^\wedge 'n$ **where**
map-vector $f \ v \equiv \chi \ i. \ f \ (v \ \$h \ i)$

definition *map-matrix* :: $('a \Rightarrow 'b) \Rightarrow 'a \ ^\wedge 'n \ ^\wedge 'm \Rightarrow 'b \ ^\wedge 'n \ ^\wedge 'm$ **where**
map-matrix $f \ A \equiv \chi \ i. \ \text{map-vector } f \ (A \ \$h \ i)$

definition *normbound* :: $'a :: \text{real-normed-field } ^\wedge 'nc \ ^\wedge 'nr \Rightarrow \text{real} \Rightarrow \text{bool}$ **where**
normbound $A \ b \equiv \forall \ x \in \text{elements-mat-h } A. \ \text{norm } x \leq b$

lemma *spectral-radius-ev-def*: $\text{spectral-radius } A = \text{Max } (\text{norm } '(\text{Collect } (\text{eigen-value } A)))$
unfolding *spectral-radius-def* *eigen-value-def* [*abs-def*]
by (*rule arg-cong* [**where** $f = \text{Max}$], *auto*)

lemma *elements-mat*: $\text{elements-mat } A = \{A \ \$$ \ (i,j) \mid i \ j. \ i < \text{dim-row } A \wedge j < \text{dim-col } A\}$
unfolding *elements-mat-def* **by** *force*

definition *vec-elements* :: $'a \ \text{Matrix.vec} \Rightarrow 'a \text{ set}$
where *vec-elements* $v = \text{set } [v \ \$ \ i. \ i <- [0 \ .. < \ \text{dim-vec } v]]$

lemma *vec-elements*: $\text{vec-elements } v = \{v \ \$ \ i \mid i. \ i < \text{dim-vec } v\}$
unfolding *vec-elements-def* **by** *auto*

context includes *vec.lifting*
begin
end

definition *from-hma_v* :: $'a \ ^\wedge 'n \Rightarrow 'a \ \text{Matrix.vec}$ **where**

$from-hma_v v = Matrix.vec\ CARD('n)\ (\lambda\ i.\ v\ \$h\ from-nat\ i)$

definition $from-hma_m :: 'a\ \hat{\ }'nc\ \hat{\ }'nr \Rightarrow 'a\ Matrix.mat$ **where**

$from-hma_m\ a = Matrix.mat\ CARD('nr)\ CARD('nc)\ (\lambda\ (i,j).\ a\ \$h\ from-nat\ i\ \$h\ from-nat\ j)$

definition $to-hma_v :: 'a\ Matrix.vec \Rightarrow 'a\ \hat{\ }'n$ **where**

$to-hma_v\ v = (\chi\ i.\ v\ \$v\ to-nat\ i)$

definition $to-hma_m :: 'a\ Matrix.mat \Rightarrow 'a\ \hat{\ }'nc\ \hat{\ }'nr$ **where**

$to-hma_m\ a = (\chi\ i\ j.\ a\ \$\$ (to-nat\ i,\ to-nat\ j))$

declare $vec-lambda-eta[simp]$

lemma $to-hma-from-hma_v[simp]: to-hma_v (from-hma_v v) = v$

by $(auto\ simp: to-hma_v-def\ from-hma_v-def\ to-nat-less-card)$

lemma $to-hma-from-hma_m[simp]: to-hma_m (from-hma_m v) = v$

by $(auto\ simp: to-hma_m-def\ from-hma_m-def\ to-nat-less-card)$

lemma $from-hma-to-hma_v[simp]:$

$v \in carrier-vec\ (CARD('n)) \Longrightarrow from-hma_v\ (to-hma_v\ v :: 'a\ \hat{\ }'n) = v$

by $(auto\ simp: to-hma_v-def\ from-hma_v-def\ to-nat-from-nat-id)$

lemma $from-hma-to-hma_m[simp]:$

$A \in carrier-mat\ (CARD('nr))\ (CARD('nc)) \Longrightarrow from-hma_m\ (to-hma_m\ A :: 'a\ \hat{\ }'nc\ \hat{\ }'nr) = A$

by $(auto\ simp: to-hma_m-def\ from-hma_m-def\ to-nat-from-nat-id)$

lemma $from-hma_v-inj[simp]: from-hma_v\ x = from-hma_v\ y \longleftrightarrow x = y$

by $(intro\ iffI,\ insert\ to-hma-from-hma_v[of\ x],\ auto)$

lemma $from-hma_m-inj[simp]: from-hma_m\ x = from-hma_m\ y \longleftrightarrow x = y$

by $(intro\ iffI,\ insert\ to-hma-from-hma_m[of\ x],\ auto)$

definition $HMA-V :: 'a\ Matrix.vec \Rightarrow 'a\ \hat{\ }'n \Rightarrow bool$ **where**

$HMA-V = (\lambda\ v\ w.\ v = from-hma_v\ w)$

definition $HMA-M :: 'a\ Matrix.mat \Rightarrow 'a\ \hat{\ }'nc\ \hat{\ }'nr \Rightarrow bool$ **where**

$HMA-M = (\lambda\ a\ b.\ a = from-hma_m\ b)$

definition $HMA-I :: nat \Rightarrow 'n :: finite \Rightarrow bool$ **where**

$HMA-I = (\lambda\ i\ a.\ i = to-nat\ a)$

context includes $lifting-syntax$

begin

lemma $Domainp-HMA-V [transfer-domain-rule]:$

$Domainp\ (HMA-V :: 'a\ Matrix.vec \Rightarrow 'a\ \hat{\ }'n \Rightarrow bool) = (\lambda\ v.\ v \in carrier-vec)$

(*CARD('n')*)
by (*intro ext iffI*, *insert from-hma-to-hma_v[symmetric]*, *auto simp: from-hma_v-def HMA-V-def*)

lemma *Domainp-HMA-M* [*transfer-domain-rule*]:
Domainp (HMA-M :: 'a Matrix.mat \Rightarrow 'a \wedge 'nc \wedge 'nr \Rightarrow bool)
*= (λ A. A \in carrier-mat *CARD('nr)* *CARD('nc)*)*
by (*intro ext iffI*, *insert from-hma-to-hma_m[symmetric]*, *auto simp: from-hma_m-def HMA-M-def*)

lemma *Domainp-HMA-I* [*transfer-domain-rule*]:
*Domainp (HMA-I :: nat \Rightarrow 'n :: finite \Rightarrow bool) = (λ i. i < *CARD('n)*) (**is** ?l = ?r)*
proof (*intro ext*)
fix i :: nat
show ?l i = ?r i
unfolding *HMA-I-def Domainp-iff*
by (*auto intro: exI[of - from-nat i] simp: to-nat-from-nat-id to-nat-less-card*)
qed

lemma *bi-unique-HMA-V* [*transfer-rule*]: *bi-unique HMA-V left-unique HMA-V right-unique HMA-V*
unfolding *HMA-V-def bi-unique-def left-unique-def right-unique-def* **by** *auto*

lemma *bi-unique-HMA-M* [*transfer-rule*]: *bi-unique HMA-M left-unique HMA-M right-unique HMA-M*
unfolding *HMA-M-def bi-unique-def left-unique-def right-unique-def* **by** *auto*

lemma *bi-unique-HMA-I* [*transfer-rule*]: *bi-unique HMA-I left-unique HMA-I right-unique HMA-I*
unfolding *HMA-I-def bi-unique-def left-unique-def right-unique-def* **by** *auto*

lemma *right-total-HMA-V* [*transfer-rule*]: *right-total HMA-V*
unfolding *HMA-V-def right-total-def* **by** *simp*

lemma *right-total-HMA-M* [*transfer-rule*]: *right-total HMA-M*
unfolding *HMA-M-def right-total-def* **by** *simp*

lemma *right-total-HMA-I* [*transfer-rule*]: *right-total HMA-I*
unfolding *HMA-I-def right-total-def* **by** *simp*

lemma *HMA-V-index* [*transfer-rule*]: (*HMA-V* \implies *HMA-I* \implies (=)) ($\$h$) ($\v)
unfolding *rel-fun-def HMA-V-def HMA-I-def from-hma_v-def*
by (*auto simp: to-nat-less-card*)

We introduce the index function to have pointwise access to HMA-matrices by a constant. Otherwise, the transfer rule with λA *i*. ($\$h$) (*A* $\$h$ *i*) instead of index is not applicable.

definition *index-hma* $A\ i\ j \equiv A\ \$h\ i\ \$h\ j$

lemma *HMA-M-index* [*transfer-rule*]:

$(HMA-M \implies HMA-I \implies HMA-I \implies (=)) (\lambda A\ i\ j. A\ \$\$ (i,j))$
index-hma

by (*intro rel-funI, simp add: index-hma-def to-nat-less-card HMA-M-def HMA-I-def from-hma_m-def*)

lemma *HMA-V-0* [*transfer-rule*]: $HMA-V\ (0_v\ CARD('n))\ (0 :: 'a :: zero \wedge 'n)$

unfolding *HMA-V-def from-hma_v-def* **by** *auto*

lemma *HMA-M-0* [*transfer-rule*]:

$HMA-M\ (0_m\ CARD('nr)\ CARD('nc))\ (0 :: 'a :: zero \wedge 'nc \wedge 'nr)$

unfolding *HMA-M-def from-hma_m-def* **by** *auto*

lemma *HMA-M-1* [*transfer-rule*]:

$HMA-M\ (1_m\ (CARD('n)))\ (mat\ 1 :: 'a :: \{zero, one\} \wedge 'n \wedge 'n)$

unfolding *HMA-M-def*

by (*auto simp add: mat-def from-hma_m-def from-nat-inj*)

lemma *from-hma_v-add*: $from-hma_v\ v + from-hma_v\ w = from-hma_v\ (v + w)$

unfolding *from-hma_v-def* **by** *auto*

lemma *HMA-V-add* [*transfer-rule*]: $(HMA-V \implies HMA-V \implies HMA-V)$

(+) (+)

unfolding *rel-fun-def HMA-V-def*

by (*auto simp: from-hma_v-add*)

lemma *from-hma_v-diff*: $from-hma_v\ v - from-hma_v\ w = from-hma_v\ (v - w)$

unfolding *from-hma_v-def* **by** *auto*

lemma *HMA-V-diff* [*transfer-rule*]: $(HMA-V \implies HMA-V \implies HMA-V)$

(-) (-)

unfolding *rel-fun-def HMA-V-def*

by (*auto simp: from-hma_v-diff*)

lemma *from-hma_m-add*: $from-hma_m\ a + from-hma_m\ b = from-hma_m\ (a + b)$

unfolding *from-hma_m-def* **by** *auto*

lemma *HMA-M-add* [*transfer-rule*]: $(HMA-M \implies HMA-M \implies HMA-M)$

(+) (+)

unfolding *rel-fun-def HMA-M-def*

by (*auto simp: from-hma_m-add*)

lemma *from-hma_m-diff*: $from-hma_m\ a - from-hma_m\ b = from-hma_m\ (a - b)$

unfolding *from-hma_m-def* **by** *auto*

lemma *HMA-M-diff* [*transfer-rule*]: $(HMA-M \implies HMA-M \implies HMA-M)$

(-) (-)

unfolding *rel-fun-def HMA-M-def*
by (*auto simp: from-hma_m-diff*)

lemma *scalar-product: fixes v :: 'a :: semiring-1 ^ 'n*
shows *scalar-prod (from-hma_v v) (from-hma_v w) = scalar-product v w*
unfolding *scalar-product-def scalar-prod-def from-hma_v-def dim-vec*
by (*simp add: sum.reindex[OF inj-to-nat, unfolded range-to-nat]*)

lemma [*simp*]:
from-hma_m (y :: 'a ^ 'nc ^ 'nr) ∈ carrier-mat (CARD('nr)) (CARD('nc))
dim-row (from-hma_m (y :: 'a ^ 'nc ^ 'nr)) = CARD('nr)
dim-col (from-hma_m (y :: 'a ^ 'nc ^ 'nr)) = CARD('nc)
unfolding *from-hma_m-def* **by** *simp-all*

lemma [*simp*]:
from-hma_v (y :: 'a ^ 'n) ∈ carrier-vec (CARD('n))
dim-vec (from-hma_v (y :: 'a ^ 'n)) = CARD('n)
unfolding *from-hma_v-def* **by** *simp-all*

declare *rel-funI* [*intro!*]

lemma *HMA-scalar-prod* [*transfer-rule*]:
(HMA-V ==> HMA-V ==> (=)) scalar-prod scalar-product
by (*auto simp: HMA-V-def scalar-product*)

lemma *HMA-row* [*transfer-rule*]: *(HMA-I ==> HMA-M ==> HMA-V) (λ i*
a. Matrix.row a i) row
unfolding *HMA-M-def HMA-I-def HMA-V-def*
by (*auto simp: from-hma_m-def from-hma_v-def to-nat-less-card row-def*)

lemma *HMA-col* [*transfer-rule*]: *(HMA-I ==> HMA-M ==> HMA-V) (λ i*
a. col a i) column
unfolding *HMA-M-def HMA-I-def HMA-V-def*
by (*auto simp: from-hma_m-def from-hma_v-def to-nat-less-card column-def*)

definition *mk-mat* :: *('i ⇒ 'j ⇒ 'c) ⇒ 'c ^ 'j ^ 'i* **where**
mk-mat f = (χ i j. f i j)

definition *mk-vec* :: *('i ⇒ 'c) ⇒ 'c ^ 'i* **where**
mk-vec f = (χ i. f i)

lemma *HMA-M-mk-mat*[*transfer-rule*]: *((HMA-I ==> HMA-I ==> (=)) ==>*
HMA-M)
(λ f. Matrix.mat (CARD('nr)) (CARD('nc)) (λ (i,j). f i j))
(mk-mat :: (('nr ⇒ 'nc ⇒ 'a) ⇒ 'a ^ 'nc ^ 'nr))

proof –

{
fix *x y i j*
assume *id: ∀ (ya :: 'nr) (yb :: 'nc). (x (to-nat ya) (to-nat yb) :: 'a) = y ya yb*

and $i: i < \text{CARD}('nr)$ **and** $j: j < \text{CARD}('nc)$
from $\text{to-nat-from-nat-id}[OF\ i]$ $\text{to-nat-from-nat-id}[OF\ j]$ $\text{id}[\text{rule-format, of from-nat } i \text{ from-nat } j]$
have $x\ i\ j = y$ $(\text{from-nat } i)$ $(\text{from-nat } j)$ **by** *auto*
}
thus *?thesis*
unfolding *rel-fun-def mk-mat-def HMA-M-def HMA-I-def from-hma_m-def* **by**
auto
qed

lemma *HMA-M-mk-vec[transfer-rule]*: $((\text{HMA-I} \implies (=)) \implies \text{HMA-V})$
 $(\lambda f. \text{Matrix.vec } (\text{CARD}('n)) (\lambda i. f\ i))$
 $(\text{mk-vec} :: (('n \Rightarrow 'a) \Rightarrow 'a \wedge 'n))$

proof –

{
fix $x\ y\ i$
assume $\text{id}: \forall (ya :: 'n). (x\ (\text{to-nat } ya) :: 'a) = y\ ya$
and $i: i < \text{CARD}('n)$
from $\text{to-nat-from-nat-id}[OF\ i]$ $\text{id}[\text{rule-format, of from-nat } i]$
have $x\ i = y$ $(\text{from-nat } i)$ **by** *auto*
}
thus *?thesis*
unfolding *rel-fun-def mk-vec-def HMA-V-def HMA-I-def from-hma_v-def* **by**
auto
qed

lemma *mat-mult-scalar*: $A ** B = \text{mk-mat } (\lambda i\ j. \text{scalar-product } (\text{row } i\ A) (\text{column } j\ B))$

unfolding *vec-eq-iff matrix-matrix-mult-def scalar-product-def mk-mat-def*
by $(\text{auto simp: row-def column-def})$

lemma *mult-mat-vec-scalar*: $A * v\ v = \text{mk-vec } (\lambda i. \text{scalar-product } (\text{row } i\ A) v)$

unfolding *vec-eq-iff matrix-vector-mult-def scalar-product-def mk-mat-def mk-vec-def*
by $(\text{auto simp: row-def column-def})$

lemma *dim-row-transfer-rule*:

$\text{HMA-M } A (A' :: 'a \wedge 'nc \wedge 'nr) \implies (=) (\text{dim-row } A) (\text{CARD}('nr))$

unfolding *HMA-M-def* **by** *auto*

lemma *dim-col-transfer-rule*:

$\text{HMA-M } A (A' :: 'a \wedge 'nc \wedge 'nr) \implies (=) (\text{dim-col } A) (\text{CARD}('nc))$

unfolding *HMA-M-def* **by** *auto*

lemma *HMA-M-mult [transfer-rule]*: $(\text{HMA-M} \implies \text{HMA-M} \implies \text{HMA-M})$

$((*) (**))$

proof –

{
fix $A\ B :: 'a :: \text{semiring-1 mat}$ **and** $A' :: 'a \wedge 'n \wedge 'nr$ **and** $B' :: 'a \wedge 'nc \wedge 'n$

```

assume 1[transfer-rule]: HMA-M A A' HMA-M B B'
note [transfer-rule] = dim-row-transfer-rule[OF 1(1)] dim-col-transfer-rule[OF
1(2)]
have HMA-M (A * B) (A' ** B')
unfolding times-mat-def mat-mult-scalar
by (transfer-prover-start, transfer-step+, transfer, auto)
}
thus ?thesis by blast
qed

```

```

lemma HMA-V-smult [transfer-rule]: ((=) ==> HMA-V ==> HMA-V) (·v)
((*s))
unfolding smult-vec-def
unfolding rel-fun-def HMA-V-def from-hma_v-def
by auto

```

```

lemma HMA-M-mult-vec [transfer-rule]: (HMA-M ==> HMA-V ==> HMA-V)
((·v)) ((*v))
proof -
{
fix A :: 'a :: semiring-1 mat and v :: 'a Matrix.vec
and A' :: 'a ^ 'nc ^ 'nr and v' :: 'a ^ 'nc
assume 1[transfer-rule]: HMA-M A A' HMA-V v v'
note [transfer-rule] = dim-row-transfer-rule
have HMA-V (A *_v v) (A' *_v v')
unfolding mult-mat-vec-def mult-mat-vec-scalar
by (transfer-prover-start, transfer-step+, transfer, auto)
}
thus ?thesis by blast
qed

```

```

lemma HMA-det [transfer-rule]: (HMA-M ==> (=)) Determinant.det
(det :: 'a :: comm-ring-1 ^ 'n ^ 'n => 'a)
proof -
{
fix a :: 'a ^ 'n ^ 'n
let ?tn = to-nat :: 'n :: finite => nat
let ?fn = from-nat :: nat => 'n
let ?zn = {0..< CARD('n)}
let ?U = UNIV :: 'n set
let ?p1 = {p. p permutes ?zn}
let ?p2 = {p. p permutes ?U}
let ?f = λ p i. if i ∈ ?U then ?fn (p (?tn i)) else i
let ?g = λ p i. ?fn (p (?tn i))
have fg: ∧ a b c. (if a ∈ ?U then b else c) = b by auto
have ?p2 = ?f ' ?p1
by (rule permutes-bij', auto simp: to-nat-less-card to-nat-from-nat-id)
hence id: ?p2 = ?g ' ?p1 by simp

```

```

have inj-g: inj-on ?g ?p1
  unfolding inj-on-def
proof (intro ballI impI ext, auto)
  fix p q i
  assume p: p permutes ?zn and q: q permutes ?zn
  and id: (λ i. ?fn (p (?tn i))) = (λ i. ?fn (q (?tn i)))
  {
    fix i
    from permutes-in-image[OF p] have pi: p (?tn i) < CARD('n) by (simp
add: to-nat-less-card)
    from permutes-in-image[OF q] have qi: q (?tn i) < CARD('n) by (simp
add: to-nat-less-card)
    from fun-cong[OF id] have ?fn (p (?tn i)) = from-nat (q (?tn i)) .
    from arg-cong[OF this, of ?tn] have p (?tn i) = q (?tn i)
    by (simp add: to-nat-from-nat-id pi qi)
  } note id = this
show p i = q i
proof (cases i < CARD('n))
  case True
  hence ?tn (?fn i) = i by (simp add: to-nat-from-nat-id)
  from id[of ?fn i, unfolded this] show ?thesis .
next
  case False
  thus ?thesis using p q unfolding permutes-def by simp
qed
qed
have mult-cong: ∧ a b c d. a = b ⇒ c = d ⇒ a * c = b * d by simp
have sum (λ p.
  signof p * (∏ i ∈ ?zn. a $h ?fn i $h ?fn (p i))) ?p1
= sum (λ p. of-int (sign p) * (∏ i ∈ UNIV. a $h i $h p i)) ?p2
  unfolding id sum.reindex[OF inj-g]
proof (rule sum.cong[OF refl], unfold mem-Collect-eq o-def, rule mult-cong)
  fix p
  assume p: p permutes ?zn
  let ?q = λ i. ?fn (p (?tn i))
  from id p have q: ?q permutes ?U by auto
  from p have pp: permutation p unfolding permutation-permutes by auto
  let ?ft = λ p i. ?fn (p (?tn i))
  have fin: finite ?zn by simp
  have sign p = sign ?q ∧ p permutes ?zn
  using p fin proof (induction rule: permutes-induct)
    case id
    show ?case by (auto simp: sign-id[unfolded id-def] permutes-id[unfolded
id-def])
  next
  case (swap a b p)
  then have ⟨permutation p⟩
  by (auto intro: permutes-imp-permutation)
  let ?sab = Transposition.transpose a b

```



```

let ?sfab = Transposition.transpose (?fn a) (?fn b)
have p-sab: permutation ?sab by (rule permutation-swap-id)
have p-sfab: permutation ?sfab by (rule permutation-swap-id)
from swap(4) have IH1: p permutes ?zn and IH2: sign p = sign (?ft p)
by auto
have sab-perm: ?sab permutes ?zn using swap(1-2) by (rule permutes-swap-id)
from permutes-compose[OF IH1 this] have perm1: ?sab o p permutes ?zn .
from IH1 have p-p1: p ∈ ?p1 by simp
hence ?ft p ∈ ?ft ' ?p1 by (rule imageI)
from this[folded id] have ?ft p permutes ?U by simp
hence p-ftp: permutation (?ft p) unfolding permutation-permutes by auto
{
  fix a b
  assume a: a ∈ ?zn and b: b ∈ ?zn
  hence (?fn a = ?fn b) = (a = b) using swap(1-2)
  by (auto simp: from-nat-inj)
} note inj = this
from inj[OF swap(1-2)] have id2: sign ?sfab = sign ?sab unfolding
sign-swap-id by simp
have id: ?ft (Transposition.transpose a b o p) = Transposition.transpose
(?fn a) (?fn b) o ?ft p
proof
  fix c
  show ?ft (Transposition.transpose a b o p) c = (Transposition.transpose
(?fn a) (?fn b) o ?ft p) c
  proof (cases p (?tn c) = a ∨ p (?tn c) = b)
    case True
    thus ?thesis by (cases, auto simp add: swap-id-eq)
  next
  case False
  hence neq: p (?tn c) ≠ a p (?tn c) ≠ b by auto
  have pc: p (?tn c) ∈ ?zn unfolding permutes-in-image[OF IH1]
  by (simp add: to-nat-less-card)
  from neq[folded inj[OF pc swap(1)] inj[OF pc swap(2)]]
  have ?fn (p (?tn c)) ≠ ?fn a ?fn (p (?tn c)) ≠ ?fn b .
  with neq show ?thesis by (auto simp: swap-id-eq)
qed
qed
show ?case unfolding IH2 id sign-compose[OF p-sab ⟨permutation p⟩]
sign-compose[OF p-sfab p-ftp] id2
by (rule conjI[OF refl perm1])
qed
thus signof p = of-int (sign ?q) unfolding sign-def by auto
show (∏ i = 0..<CARD('n). a $h ?fn i $h ?fn (p i)) =
(∏ i ∈ UNIV. a $h i $h ?q i) unfolding
range-to-nat[symmetric] prod.reindex[OF inj-to-nat]
by (rule prod.cong[OF refl], unfold o-def, simp)
qed
}

```

thus *?thesis* **unfolding** *HMA-M-def*
by (*auto simp: from-hma_m-def Determinant.det-def det-def*)
qed

lemma *HMA-mat[transfer-rule]*: ((=) ==> *HMA-M*) ($\lambda k. k \cdot_m 1_m \text{ CARD}('n)$)

(*Finite-Cartesian-Product.mat* :: '*a*::*semiring-1* \Rightarrow '*a*^{*n*}^{*n*})
unfolding *Finite-Cartesian-Product.mat-def[abs-def] rel-fun-def HMA-M-def*
by (*auto simp: from-hma_m-def from-nat-inj*)

lemma *HMA-mat-minus[transfer-rule]*: (*HMA-M* ==> *HMA-M* ==> *HMA-M*)

($\lambda A B. A + \text{map-mat } \text{uminus } B$) ((-)) :: '*a* :: *group-add* $\hat{\ }^n \hat{\ }^n \Rightarrow$ '*a*^{*n*}^{*n*}
 \Rightarrow '*a*^{*n*}^{*n*})
unfolding *rel-fun-def HMA-M-def from-hma_m-def* **by** *auto*

definition *mat2matofpoly* **where** *mat2matofpoly* *A* = ($\chi i j. [: A \$ i \$ j :]$)

definition *charpoly* **where** *charpoly-def*: *charpoly* *A* = *det* (*mat* (*monom* 1 (*Suc* 0)) - *mat2matofpoly* *A*)

definition *erase-mat* :: '*a* :: *zero* $\hat{\ }^n \hat{\ }^n \Rightarrow$ '*nr* \Rightarrow '*nc* \Rightarrow '*a* $\hat{\ }^n \hat{\ }^n$
where *erase-mat* *A* *i j* = ($\chi i'. \chi j'. \text{if } i' = i \vee j' = j \text{ then } 0 \text{ else } A \$ i' \$ j'$)

definition *sum-UNIV-type* :: ('*n* :: *finite* \Rightarrow '*a* :: *comm-monoid-add*) \Rightarrow '*n* *itself*
 \Rightarrow '*a* **where**
sum-UNIV-type *f* - = *sum* *f* *UNIV*

definition *sum-UNIV-set* :: (*nat* \Rightarrow '*a* :: *comm-monoid-add*) \Rightarrow *nat* \Rightarrow '*a* **where**
sum-UNIV-set *f* *n* = *sum* *f* {..*n*}

definition *HMA-T* :: *nat* \Rightarrow '*n* :: *finite* *itself* \Rightarrow *bool* **where**
HMA-T *n* - = (*n* = *CARD*('*n*))

lemma *HMA-mat2matofpoly[transfer-rule]*: (*HMA-M* ==> *HMA-M*) ($\lambda x. \text{map-mat}$
($\lambda a. [: a :]$) *x*) *mat2matofpoly*
unfolding *rel-fun-def HMA-M-def from-hma_m-def mat2matofpoly-def* **by** *auto*

lemma *HMA-char-poly* [*transfer-rule*]:
((*HMA-M* :: ('*a*:: *comm-ring-1* *mat* \Rightarrow '*a*^{*n*}^{*n*} \Rightarrow *bool*)) ==> (=)) *char-poly*
charpoly

proof -

{
fix *A* :: '*a* *mat* **and** *A'* :: '*a*^{*n*}^{*n*}
assume [*transfer-rule*]: *HMA-M* *A* *A'*
hence [*simp*]: *dim-row* *A* = *CARD*('*n*) **by** (*simp* *add: HMA-M-def*)
have [*simp*]: *monom* 1 (*Suc* 0) = [: 0, 1 :: '*a* :]
by (*simp* *add: monom-Suc*)

```

have [simp]: map-mat uminus (map-mat ( $\lambda a. [:a:]$ ) A) = map-mat ( $\lambda a. [-a:]$ )
A
  by (rule eq-matI, auto)
have char-poly A = charpoly A'
  unfolding char-poly-def[abs-def] char-poly-matrix-def charpoly-def[abs-def]
  by (transfer, simp)
}
thus ?thesis by blast
qed

```

```

lemma HMA-eigen-vector [transfer-rule]: (HMA-M ==> HMA-V ==> (=))
eigenvector eigen-vector
proof -
{
  fix A :: 'a mat and v :: 'a Matrix.vec
  and A' :: 'a ^ 'n ^ 'n and v' :: 'a ^ 'n and k :: 'a
  assume 1[transfer-rule]: HMA-M A A' and 2[transfer-rule]: HMA-V v v'
  hence [simp]: dim-row A = CARD('n) dim-vec v = CARD('n) by (auto simp
add: HMA-V-def HMA-M-def)
  have [simp]: v  $\in$  carrier-vec CARD('n) using 2 unfolding HMA-V-def by
simp
  have eigenvector A v = eigen-vector A' v'
  unfolding eigenvector-def[abs-def] eigen-vector-def[abs-def]
  by (transfer, simp)
}
thus ?thesis by blast
qed

```

```

lemma HMA-eigen-value [transfer-rule]: (HMA-M ==> (=) ==> (=)) eigen-
value eigen-value
proof -
{
  fix A :: 'a mat and A' :: 'a ^ 'n ^ 'n and k
  assume 1[transfer-rule]: HMA-M A A'
  hence [simp]: dim-row A = CARD('n) by (simp add: HMA-M-def)
  note [transfer-rule] = dim-row-transfer-rule[OF 1(1)]
  have (eigenvalue A k) = (eigen-value A' k)
  unfolding eigenvalue-def[abs-def] eigen-value-def[abs-def]
  by (transfer, auto simp add: eigenvector-def)
}
thus ?thesis by blast
qed

```

```

lemma HMA-spectral-radius [transfer-rule]:
(HMA-M ==> (=)) Spectral-Radius.spectral-radius spectral-radius
unfolding Spectral-Radius.spectral-radius-def[abs-def] spectrum-def

```

```

    spectral-radius-ev-def[abs-def]
  by transfer-prover

lemma HMA-elements-mat[transfer-rule]: ((HMA-M :: ('a mat  $\Rightarrow$  'a  $\wedge$  'nc  $\wedge$  'nr
 $\Rightarrow$  bool))  $\implies$  (=))
  elements-mat elements-mat-h
proof -
{
  fix y :: 'a  $\wedge$  'nc  $\wedge$  'nr and i j :: nat
  assume i: i < CARD('nr) and j: j < CARD('nc)
  hence from-hmam y $$ (i, j)  $\in$  range ( $\lambda$ (i, ya). y $h i $h ya)
    using to-nat-from-nat-id[OF i] to-nat-from-nat-id[OF j] by (auto simp:
from-hmam-def)
}
moreover
{
  fix y :: 'a  $\wedge$  'nc  $\wedge$  'nr and a b
  have  $\exists$  i j. y $h a $h b = from-hmam y $$ (i, j)  $\wedge$  i < CARD('nr)  $\wedge$  j <
CARD('nc)
  unfolding from-hmam-def
  by (rule exI[of - Bij-Nat.to-nat a], rule exI[of - Bij-Nat.to-nat b], auto
simp: to-nat-less-card)
}
ultimately show ?thesis
  unfolding elements-mat[abs-def] elements-mat-h-def[abs-def] HMA-M-def
  by auto
qed

lemma HMA-vec-elements[transfer-rule]: ((HMA-V :: ('a Matrix.vec  $\Rightarrow$  'a  $\wedge$  'n  $\Rightarrow$ 
bool))  $\implies$  (=))
  vec-elements vec-elements-h
proof -
{
  fix y :: 'a  $\wedge$  'n and i :: nat
  assume i: i < CARD('n)
  hence from-hmav y $ i  $\in$  range (vec-nth y)
    using to-nat-from-nat-id[OF i] by (auto simp: from-hmav-def)
}
moreover
{
  fix y :: 'a  $\wedge$  'n and a
  have  $\exists$  i. y $h a = from-hmav y $ i  $\wedge$  i < CARD('n)
  unfolding from-hmav-def
  by (rule exI[of - Bij-Nat.to-nat a], auto simp: to-nat-less-card)
}
ultimately show ?thesis
  unfolding vec-elements[abs-def] vec-elements-h-def[abs-def] rel-fun-def HMA-V-def
  by auto
qed

```

lemma *norm-bound-elements-mat*: *norm-bound* A $b = (\forall x \in \text{elements-mat } A. \text{norm } x \leq b)$

unfolding *norm-bound-def elements-mat* **by** *auto*

lemma *HMA-normbound* [*transfer-rule*]:

$((\text{HMA-M} :: 'a :: \text{real-normed-field mat} \Rightarrow 'a \wedge 'nc \wedge 'nr \Rightarrow \text{bool}) \implies (=) \implies (=))$

norm-bound normbound

unfolding *normbound-def[abs-def] norm-bound-elements-mat[abs-def]*

by (*transfer-prover*)

lemma *HMA-map-matrix* [*transfer-rule*]:

$((=) \implies \text{HMA-M} \implies \text{HMA-M}) \text{ map-mat map-matrix}$

unfolding *map-vector-def map-matrix-def[abs-def] map-mat-def[abs-def] HMA-M-def from-hma_m-def*

by *auto*

lemma *HMA-transpose-matrix* [*transfer-rule*]:

$(\text{HMA-M} \implies \text{HMA-M}) \text{ transpose-mat transpose}$

unfolding *transpose-mat-def transpose-def HMA-M-def from-hma_m-def* **by** *auto*

lemma *HMA-map-vector* [*transfer-rule*]:

$((=) \implies \text{HMA-V} \implies \text{HMA-V}) \text{ map-vec map-vector}$

unfolding *map-vector-def[abs-def] map-vec-def[abs-def] HMA-V-def from-hma_v-def*

by *auto*

lemma *HMA-similar-mat-wit* [*transfer-rule*]:

$((\text{HMA-M} :: - \Rightarrow 'a :: \text{comm-ring-1} \wedge 'n \wedge 'n \Rightarrow -) \implies \text{HMA-M} \implies \text{HMA-M} \implies \text{HMA-M} \implies (=))$

similar-mat-wit similar-matrix-wit

proof (*intro rel-funI, goal-cases*)

case ($1 a A b B c C d D$)

note [*transfer-rule*] = *this*

hence *id*: $\text{dim-row } a = \text{CARD}(n)$ **by** (*auto simp: HMA-M-def*)

have $*$: $(c * d = 1_m (\text{dim-row } a) \wedge d * c = 1_m (\text{dim-row } a) \wedge a = c * b * d = (C ** D = \text{mat } 1 \wedge D ** C = \text{mat } 1 \wedge A = C ** B ** D))$ **unfolding** *id*

by (*transfer, simp*)

show $?case$ **unfolding** *similar-mat-wit-def Let-def similar-matrix-wit-def **

using 1 **by** (*auto simp: HMA-M-def*)

qed

lemma *HMA-similar-mat* [*transfer-rule*]:

$((\text{HMA-M} :: - \Rightarrow 'a :: \text{comm-ring-1} \wedge 'n \wedge 'n \Rightarrow -) \implies \text{HMA-M} \implies (=))$

similar-mat similar-matrix

proof (*intro rel-funI, goal-cases*)

case ($1 a A b B$)

note [*transfer-rule*] = *this*

hence $id: dim\text{-}row\ a = CARD('n)$ **by** (*auto simp: HMA-M-def*)
{
 fix $c\ d$
 assume *similar-mat-wit a b c d*
 hence $\{c,d\} \subseteq carrier\text{-}mat\ CARD('n)$ $CARD('n)$ **unfolding** *similar-mat-wit-def*
id Let-def **by** *auto*
} **note** $*$ = *this*
show *?case* **unfolding** *similar-mat-def similar-matrix-def*
 by (*transfer, insert *, blast*)
qed

lemma *HMA-spectrum[transfer-rule]*: ($HMA\text{-}M \implies (=)$) *spectrum Spectrum*
 unfolding *spectrum-def[abs-def] Spectrum-def[abs-def]*
 by *transfer-prover*

lemma *HMA-M-erase-mat[transfer-rule]*: ($HMA\text{-}M \implies HMA\text{-}I \implies HMA\text{-}I \implies HMA\text{-}M$) *mat-erase erase-mat*
 unfolding *mat-erase-def[abs-def] erase-mat-def[abs-def]*
 by (*auto simp: HMA-M-def HMA-I-def from-hma_m-def to-nat-from-nat-id intro!: eq-matI*)

lemma *HMA-M-sum-UNIV[transfer-rule]*:
 ($HMA\text{-}I \implies (=) \implies HMA\text{-}T \implies (=)$) *sum-UNIV-set sum-UNIV-type*
 unfolding *rel-fun-def*
proof (*clarify, rename-tac f fT n nT*)
 fix f **and** $fT :: 'b \Rightarrow 'a$ **and** n **and** $nT :: 'b$ *itself*
 assume $f: \forall x\ y. HMA\text{-}I\ x\ y \longrightarrow f\ x = fT\ y$
 and $n: HMA\text{-}T\ n\ nT$
 let $?f = from\text{-}nat :: nat \Rightarrow 'b$
 let $?t = to\text{-}nat :: 'b \Rightarrow nat$
 from $n[unfolded\ HMA\text{-}T\text{-}def]$ **have** $n: n = CARD('b)$.
 from *to-nat-from-nat-id* **where** $'a = 'b$, *folded n*
 have $tf: i < n \implies ?t\ (?f\ i) = i$ **for** i **by** *auto*
 have $sum\text{-}UNIV\text{-}set\ f\ n = sum\ f\ (?t\ ' ?f\ \{..\ < n\})$
 unfolding *sum-UNIV-set-def*
 by (*rule arg-cong[of - - sum f], insert tf, force*)
 also **have** $\dots = sum\ (f \circ ?t)\ (?f\ \{..\ < n\})$
 by (*rule sum.reindex, insert tf n, auto simp: inj-on-def*)
 also **have** $?f\ \{..\ < n\} = UNIV$
 using *range-to-nat* **where** $'a = 'b$, *folded n* **by** *force*
 also **have** $sum\ (f \circ ?t)\ UNIV = sum\ fT\ UNIV$
 proof (*rule sum.cong[OF refl]*)
 fix $i :: 'b$
 show $(f \circ ?t)\ i = fT\ i$ **unfolding** *o-def*
 by (*rule f[rule-format], auto simp: HMA-I-def*)
qed
 also **have** $\dots = sum\text{-}UNIV\text{-}type\ fT\ nT$
 unfolding *sum-UNIV-type-def ..*
 finally **show** $sum\text{-}UNIV\text{-}set\ f\ n = sum\text{-}UNIV\text{-}type\ fT\ nT$.

qed
end

Setup a method to easily convert theorems from JNF into HMA.

method *transfer-hma* **uses** *rule* = (
 (*fold index-hma-def*)?,
transfer,
rule rule,
 (*unfold carrier-vec-def carrier-mat-def*)?,
auto)

Now it becomes easy to transfer results which are not yet proven in HMA, such as:

lemma *matrix-add-vect-distrib*: $(A + B) * v v = A * v v + B * v v$
by (*transfer-hma rule: add-mult-distrib-mat-vec*)

lemma *matrix-vector-right-distrib*: $M * v (v + w) = M * v v + M * v w$
by (*transfer-hma rule: mult-add-distrib-mat-vec*)

lemma *matrix-vector-right-distrib-diff*: $(M :: 'a :: \text{ring-1 } ^n \text{ } ^n) * v (v - w) = M * v v - M * v w$
by (*transfer-hma rule: mult-minus-distrib-mat-vec*)

lemma *eigen-value-root-charpoly*:
eigen-value $A k \longleftrightarrow \text{poly } (\text{charpoly } (A :: 'a :: \text{field } ^n \text{ } ^n)) k = 0$
by (*transfer-hma rule: eigenvalue-root-char-poly*)

lemma *finite-spectrum*: **fixes** $A :: 'a :: \text{field } ^n \text{ } ^n$
shows *finite* (*Collect* (*eigen-value* A))
by (*transfer-hma rule: card-finite-spectrum(1)[unfolded spectrum-def]*)

lemma *non-empty-spectrum*: **fixes** $A :: \text{complex } ^n \text{ } ^n$
shows *Collect* (*eigen-value* A) $\neq \{\}$
by (*transfer-hma rule: spectrum-non-empty[unfolded spectrum-def]*)

lemma *charpoly-transpose*: $\text{charpoly } (\text{transpose } A :: 'a :: \text{field } ^n \text{ } ^n) = \text{charpoly } A$
by (*transfer-hma rule: char-poly-transpose-mat*)

lemma *eigen-value-transpose*: $\text{eigen-value } (\text{transpose } A :: 'a :: \text{field } ^n \text{ } ^n) v = \text{eigen-value } A v$
unfolding *eigen-value-root-charpoly charpoly-transpose* **by** *simp*

lemma *matrix-diff-vect-distrib*: $(A - B) * v v = A * v v - B * v v$ ($v :: 'a :: \text{ring-1 } ^n \text{ } ^n$)
by (*transfer-hma rule: minus-mult-distrib-mat-vec*)

lemma *similar-matrix-charpoly*: $\text{similar-matrix } A B \implies \text{charpoly } A = \text{charpoly } B$

by (transfer-hma rule: char-poly-similar)

lemma *pderiv-char-poly-erase-mat*: fixes $A :: 'a :: idom \hat{\ } ^ 'n \hat{\ } ^ 'n$
 shows $monom\ 1\ 1 * pderiv\ (charpoly\ A) = sum\ (\lambda\ i.\ charpoly\ (erase-mat\ A\ i\ i))\ UNIV$
proof –
 let $?A = from-hma_m\ A$
 let $?n = CARD('n)$
 have $tA[transfer-rule]: HMA-M\ ?A\ A$ unfolding *HMA-M-def* by *simp*
 have $tN[transfer-rule]: HMA-T\ ?n\ TYPE('n)$ unfolding *HMA-T-def* by *simp*
 have $A: ?A \in carrier-mat\ ?n\ ?n$ unfolding *from-hma_m-def* by *auto*
 have $id: sum\ (\lambda\ i.\ charpoly\ (erase-mat\ A\ i\ i))\ UNIV =$
 $sum-UNIV-type\ (\lambda\ i.\ charpoly\ (erase-mat\ A\ i\ i))\ TYPE('n)$
 unfolding *sum-UNIV-type-def* ..
 show *?thesis* unfolding *id*
 by (transfer, insert *pderiv-char-poly-mat-erase[OF A]*, *simp add: sum-UNIV-set-def*)
qed

lemma *degree-monic-charpoly*: fixes $A :: 'a :: comm-ring-1 \hat{\ } ^ 'n \hat{\ } ^ 'n$
 shows $degree\ (charpoly\ A) = CARD('n) \wedge monic\ (charpoly\ A)$
proof (*transfer, goal-cases*)
 case 1
 from *degree-monic-char-poly[OF 1]* show *?case* by *auto*
qed

end

4 Perron-Frobenius Theorem

4.1 Auxiliary Notions

We define notions like non-negative real-valued matrix, both in JNF and in HMA. These notions will be linked via HMA-connect.

theory *Perron-Frobenius-Aux*
imports *HMA-Connect*
begin

definition *real-nonneg-mat* :: $complex\ mat \Rightarrow bool$ **where**
real-nonneg-mat $A \equiv \forall\ a \in elements-mat\ A.\ a \in \mathbb{R} \wedge Re\ a \geq 0$

definition *real-nonneg-vec* :: $complex\ Matrix.vec \Rightarrow bool$ **where**
real-nonneg-vec $v \equiv \forall\ a \in vec-elements\ v.\ a \in \mathbb{R} \wedge Re\ a \geq 0$

definition *real-non-neg-vec* :: $complex \hat{\ } ^ 'n \Rightarrow bool$ **where**
real-non-neg-vec $v \equiv (\forall\ a \in vec-elements-h\ v.\ a \in \mathbb{R} \wedge Re\ a \geq 0)$

definition *real-non-neg-mat* :: $complex \hat{\ } ^ 'nr \hat{\ } ^ 'nc \Rightarrow bool$ **where**
real-non-neg-mat $A \equiv (\forall\ a \in elements-mat-h\ A.\ a \in \mathbb{R} \wedge Re\ a \geq 0)$

lemma *real-non-neg-matD*: **assumes** *real-non-neg-mat A*
shows $A \ \$h \ i \ \$h \ j \in \mathbf{R} \ Re \ (A \ \$h \ i \ \$h \ j) \geq 0$
using *assms unfolding real-non-neg-mat-def elements-mat-h-def* **by** *auto*

definition *nonneg-mat* :: '*a* :: *linordered-idom mat* \Rightarrow *bool* **where**
nonneg-mat A $\equiv \forall a \in \text{elements-mat } A. a \geq 0$

definition *non-neg-mat* :: '*a* :: *linordered-idom* \wedge '*nr* \wedge '*nc* \Rightarrow *bool* **where**
non-neg-mat A $\equiv (\forall a \in \text{elements-mat-h } A. a \geq 0)$

context includes *lifting-syntax*
begin

lemma *HMA-real-non-neg-mat* [*transfer-rule*]:
 $((HMA-M :: \text{complex mat} \Rightarrow \text{complex} \wedge 'nc \wedge 'nr \Rightarrow \text{bool}) \implies (=))$
real-nonneg-mat real-non-neg-mat
unfolding *real-nonneg-mat-def[abs-def] real-non-neg-mat-def[abs-def]*
by *transfer-prover*

lemma *HMA-real-non-neg-vec* [*transfer-rule*]:
 $((HMA-V :: \text{complex Matrix.vec} \Rightarrow \text{complex} \wedge 'n \Rightarrow \text{bool}) \implies (=))$
real-nonneg-vec real-non-neg-vec
unfolding *real-nonneg-vec-def[abs-def] real-non-neg-vec-def[abs-def]*
by *transfer-prover*

lemma *HMA-non-neg-mat* [*transfer-rule*]:
 $((HMA-M :: 'a :: \text{linordered-idom mat} \Rightarrow 'a \wedge 'nc \wedge 'nr \Rightarrow \text{bool}) \implies (=))$
nonneg-mat non-neg-mat
unfolding *nonneg-mat-def[abs-def] non-neg-mat-def[abs-def]*
by *transfer-prover*

end

primrec *matpow* :: '*a*::*semiring-1* \wedge '*n* \wedge '*n* \Rightarrow *nat* \Rightarrow '*a* \wedge '*n* \wedge '*n* **where**
matpow-0: *matpow A 0 = mat 1 |*
matpow-Suc: *matpow A (Suc n) = (matpow A n) ** A*

context includes *lifting-syntax*
begin

lemma *HMA-pow-mat*[*transfer-rule*]:
 $((HMA-M :: 'a :: \{\text{semiring-1}\} \text{ mat} \Rightarrow 'a \wedge 'n \wedge 'n \Rightarrow \text{bool}) \implies (=) \implies HMA-M)$
pow-mat matpow

proof –

{
fix *A* :: '*a* *mat* **and** *A'* :: '*a* \wedge '*n* \wedge '*n* **and** *n* :: *nat*
assume [*transfer-rule*]: *HMA-M A A'*
hence [*simp*]: *dim-row A = CARD('n)* **unfolding** *HMA-M-def* **by** *simp*

```

    have HMA-M (pow-mat A n) (matpow A' n)
    proof (induct n)
      case (Suc n)
      note [transfer-rule] = this
      show ?case by (simp, transfer-prover)
    qed (simp, transfer-prover)
  }
  thus ?thesis by blast
qed
end

```

```

lemma trancl-image:
   $(i,j) \in R^+ \implies (f i, f j) \in (\text{map-prod } f f \text{ ' } R)^+$ 
proof (induct rule: trancl-induct)
  case (step j k)
  from step(2) have  $(f j, f k) \in \text{map-prod } f f \text{ ' } R$  by auto
  from step(3) this show ?case by auto
qed auto

```

```

lemma inj-trancl-image: assumes inj: inj f
  shows  $(f i, f j) \in (\text{map-prod } f f \text{ ' } R)^+ = ((i,j) \in R^+) \text{ (is ?l = ?r)}$ 
proof
  assume ?r from trancl-image[OF this] show ?l .
next
  assume ?l from trancl-image[OF this, of the-inv f]
  show ?r unfolding image-image prod.map-comp o-def the-inv-f-f[OF inj] by
  auto
qed

```

```

lemma matrix-add-rdistrib:  $((B + C) ** A) = (B ** A) + (C ** A)$ 
  by (vector matrix-matrix-mult-def sum.distrib[symmetric] field-simps)

```

```

lemma norm-smult:  $\text{norm } ((a :: \text{real}) *s x) = \text{abs } a * \text{norm } x$ 
  unfolding norm-vec-def
  by (metis norm-scaleR norm-vec-def scalar-mult-eq-scaleR)

```

```

lemma nonneg-mat-mult:
   $\text{nonneg-mat } A \implies \text{nonneg-mat } B \implies A \in \text{carrier-mat } nr \ n$ 
 $\implies B \in \text{carrier-mat } n \ nc \implies \text{nonneg-mat } (A * B)$ 
  unfolding nonneg-mat-def
  by (auto simp: elements-mat-def scalar-prod-def intro!: sum-nonneg)

```

```

lemma nonneg-mat-power: assumes  $A \in \text{carrier-mat } n \ n \ \text{nonneg-mat } A$ 
  shows  $\text{nonneg-mat } (A \hat{\ }_m \ k)$ 
proof (induct k)
  case 0
  thus ?case by (auto simp: nonneg-mat-def)
next
  case (Suc k)

```

from *nonneg-mat-mult*[*OF this assms(2) - assms(1), of n*] *assms(1)*
show *?case by auto*
qed

lemma *nonneg-matD*: **assumes** *nonneg-mat A*
and $i < \text{dim-row } A$ **and** $j < \text{dim-col } A$
shows $A \text{ \$(\$ (i,j) \ge 0)}$
using *assms unfolding nonneg-mat-def elements-mat by auto*

lemma (**in** *comm-ring-hom*) *similar-mat-wit-hom*: **assumes**
similar-mat-wit A B C D
shows *similar-mat-wit (mat_h A) (mat_h B) (mat_h C) (mat_h D)*
proof –

obtain n **where** $n: n = \text{dim-row } A$ **by auto**
note $*$ = *similar-mat-witD*[*OF n assms*]
from $*$ **have** [*simp*]: $\text{dim-row } C = n$ **by auto**
note $C = *(6)$ **note** $D = *(7)$
note $id = \text{mat-hom-mult}$ [*OF C D*] *mat-hom-mult*[*OF D C*]
note $** = *(1-3)$ [*THEN arg-cong*[*of - - mat_h*], *unfolded id*]
note $mult = \text{mult-carrier-mat}$ [*of - n n*]
note $\text{hom-mult} = \text{mat-hom-mult}$ [*of - n n - n*]
show *?thesis unfolding similar-mat-wit-def Let-def unfolding ** (3) using*
 $** (1,2)$
by (*auto simp: n[symmetric] hom-mult simp: *(4-) mult*)
qed

lemma (**in** *comm-ring-hom*) *similar-mat-hom*:
similar-mat A B \implies similar-mat (mat_h A) (mat_h B)
using *similar-mat-wit-hom*[*of A B C D for C D*]
by (*smt similar-mat-def*)

lemma *det-dim-1*: **assumes** $A: A \in \text{carrier-mat } n \ n$
and $n: n = 1$
shows *Determinant.det A = A \\$(\\$ (0,0)*
by (*subst laplace-expansion-column*[*OF A[unfolded n], of 0*], *insert A n*,
auto simp: cofactor-def mat-delete-def)

lemma *det-dim-2*: **assumes** $A: A \in \text{carrier-mat } n \ n$
and $n: n = 2$
shows *Determinant.det A = A \\$(\\$ (0,0) * A \\$(\\$ (1,1) - A \\$(\\$ (0,1) * A \\$(\\$ (1,0)*
proof –
have *set: ($\sum i < (2 :: nat). f i) = f 0 + f 1$ for f*
by (*subst sum.cong*[*of - {0,1} f f*], *auto*)
show *?thesis*
apply (*subst laplace-expansion-column*[*OF A[unfolded n], of 0*], *insert A n*,
auto simp: cofactor-def mat-delete-def set)
apply (*subst (1 2) det-dim-1, auto*)
done
qed

lemma *jordan-nf-root-char-poly*: **fixes** $A :: 'a :: \{\text{semiring-no-zero-divisors, idom}\}$
mat
assumes *jordan-nf* A *n-as*
and $(m, lam) \in \text{set } n\text{-as}$
shows $\text{poly } (\text{char-poly } A) \text{ lam} = 0$
proof –
from *assms* **have** $m0: m \neq 0$ **unfolding** *jordan-nf-def* **by** *force*
from *split-list[OF assms(2)]* **obtain** *as bs* **where** $nas: n\text{-as} = as @ (m, lam) \#$
bs **by** *auto*
show *?thesis* **using** $m0$
unfolding *jordan-nf-char-poly[OF assms(1)]* *nas poly-prod-list prod-list-zero-iff*
by $(\text{auto simp: } o\text{-def})$
qed

lemma *inverse-power-tendsto-zero*:
 $(\lambda x. \text{inverse } ((\text{of-nat } x :: 'a :: \text{real-normed-div-algebra}) \wedge \text{Suc } d)) \longrightarrow 0$
proof $(\text{rule filterlim-compose[OF tendsto-inverse-0]},$
*intro filterlim-at-infinity[THEN iffD2, of 0] allI impI, goal-cases)
case $(2\ r)$
let $?r = \text{nat } (\text{ceiling } r) + 1$
show *?case*
proof $(\text{intro eventually-sequentiallyI[of ?r], unfold norm-power norm-of-nat})$
fix x
assume $r: ?r \leq x$
hence $x1: \text{real } x \geq 1$ **by** *auto*
have $r \leq \text{real } ?r$ **by** *linarith*
also have $\dots \leq x$ **using** r **by** *auto*
also have $\dots \leq \text{real } x \wedge \text{Suc } d$ **using** $x1$ **by** *simp*
finally show $r \leq \text{real } x \wedge \text{Suc } d$.
qed
qed *simp**

lemma *inverse-of-nat-tendsto-zero*:
 $(\lambda x. \text{inverse } (\text{of-nat } x :: 'a :: \text{real-normed-div-algebra})) \longrightarrow 0$
using *inverse-power-tendsto-zero[of 0]* **by** *auto*

lemma *poly-times-exp-tendsto-zero*: **assumes** $b: \text{norm } (b :: 'a :: \text{real-normed-field})$
 < 1
shows $(\lambda x. \text{of-nat } x \wedge k * b \wedge x) \longrightarrow 0$
proof $(\text{cases } b = 0)$
case *False*
define nla **where** $nla = \text{norm } b$
define s **where** $s = \text{sqrt } nla$
from b *False* **have** $nla: 0 < nla \wedge nla < 1$ **unfolding** *nla-def* **by** *auto*
hence $s: 0 < s \wedge s < 1$ **unfolding** *s-def* **by** *auto*
{
fix x

```

have  $s^{\wedge}x * s^{\wedge}x = \text{sqrt} (nla^{\wedge}(2 * x))$ 
  unfolding s-def power-add[symmetric]
  unfolding real-sqrt-power[symmetric]
  by (rule arg-cong[of - -  $\lambda x. \text{sqrt} (nla^{\wedge}x)$ ], simp)
also have  $\dots = nla^{\wedge}x$  unfolding power-mult real-sqrt-power
  using nla by simp
finally have  $nla^{\wedge}x = s^{\wedge}x * s^{\wedge}x$  by simp
} note nla-s = this
show ( $\lambda x. \text{of-nat } x^{\wedge}k * b^{\wedge}x \longrightarrow 0$ )
proof (rule tendsto-norm-zero-cancel, unfold norm-mult norm-power norm-of-nat
nla-def[symmetric] nla-s
  mult.assoc[symmetric])
  from poly-exp-constant-bound[OF s, of 1 k] obtain p where
    p: real  $x^{\wedge}k * s^{\wedge}x \leq p$  for  $x$  by (auto simp: ac-simps)
  have norm (real  $x^{\wedge}k * s^{\wedge}x) = real  $x^{\wedge}k * s^{\wedge}x$  for  $x$  using  $s$  by auto$ 
  with p have p: norm (real  $x^{\wedge}k * s^{\wedge}x) \leq p$  for  $x$  by auto
  from s have s: norm  $s < 1$  by auto
  show ( $\lambda x. real  $x^{\wedge}k * s^{\wedge}x * s^{\wedge}x \longrightarrow 0$$ )
    by (rule lim-null-mult-left-bounded[OF - LIMSEQ-power-zero[OF s], of - p],
insert p, auto)
  qed
next
  case True
  show ?thesis unfolding True
    by (subst tendsto-cong[of -  $\lambda x. 0$ ], rule eventually-sequentiallyI[of 1], auto)
  qed

lemma (in linorder-topology) tendsto-Min: assumes  $I: I \neq \{\}$  and fin: finite I
  shows ( $\bigwedge i. i \in I \implies (f i \longrightarrow a i) F \implies ((\lambda x. \text{Min} ((\lambda i. f i x) ' I)) \longrightarrow$ 
  (Min (a ' I) :: 'a)) F)
  using fin I
proof (induct rule: finite-induct)
  case (insert i I)
  hence i: (f i  $\longrightarrow$  a i) F by auto
  show ?case
  proof (cases I = \{\})
    case True
    show ?thesis unfolding True using i by auto
  next
  case False
  have *: Min (a ' insert i I) = min (a i) (Min (a ' I)) using False insert(1)
by auto
  have **: ( $\lambda x. \text{Min} ((\lambda i. f i x) ' \text{insert } i I) = (\lambda x. \text{min} (f i x) (\text{Min} ((\lambda i. f i x) ' I)))$ )
    using False insert(1) by auto
  have IH: ( $(\lambda x. \text{Min} ((\lambda i. f i x) ' I)) \longrightarrow \text{Min} (a ' I) F$ )
    using insert(3)[OF insert(4) False] by auto
  show ?thesis unfolding * **

```

by (auto intro!: tendsto-min i IH)
qed
qed simp

lemma tendsto-mat-mult [tendsto-intros]:
 $(f \longrightarrow a) F \Longrightarrow (g \longrightarrow b) F \Longrightarrow ((\lambda x. f x ** g x) \longrightarrow a ** b) F$
for $f :: 'a \Rightarrow 'b :: \{\text{semiring-1}, \text{real-normed-algebra}\}^{\wedge 'n1} \wedge 'n2$
unfolding matrix-matrix-mult-def[abs-def] **by** (auto intro!: tendsto-intros)

lemma tendsto-matpower [tendsto-intros]: $(f \longrightarrow a) F \Longrightarrow ((\lambda x. \text{matpow } (f x) n) \longrightarrow \text{matpow } a n) F$
for $f :: 'a \Rightarrow 'b :: \{\text{semiring-1}, \text{real-normed-algebra}\}^{\wedge 'n} \wedge 'n$
by (induct n, simp-all add: tendsto-mat-mult)

lemma continuous-matpow: continuous-on $R (\lambda A :: 'a :: \{\text{semiring-1}, \text{real-normed-algebra-1}\}^{\wedge 'n} \wedge 'n. \text{matpow } A n)$
unfolding continuous-on-def **by** (auto intro!: tendsto-intros)

lemma vector-smult-distrib: $(A *v ((a :: 'a :: \text{comm-ring-1}) *s x)) = a *s ((A *v x))$
unfolding matrix-vector-mult-def vector-scalar-mult-def
by (simp add: ac-simps sum-distrib-left)

instance real :: ordered-semiring-strict
by (intro-classes, auto)

lemma poly-tendsto-pinfty: **fixes** $p :: \text{real poly}$
assumes lead-coeff $p > 0$ degree $p \neq 0$
shows $\text{poly } p \longrightarrow \infty$
unfolding Lim-PInfy
proof
fix b
show $\exists N. \forall n \geq N. \text{ereal } b \leq \text{ereal } (\text{poly } p (\text{real } n))$
unfolding erreal-less-eq **using** poly-pinfty-ge[OF assms, of b]
by (meson of-nat-le-iff order-trans real-arch-simple)
qed

lemma div-lt-nat: $(j :: \text{nat}) < x * y \Longrightarrow j \text{ div } x < y$
by (simp add: less-mult-imp-div-less mult.commute)

definition diagvector :: $('n \Rightarrow 'a :: \text{semiring-0}) \Rightarrow 'a^{\wedge 'n} \wedge 'n$ **where**
diagvector $x = (\chi i. \chi j. \text{if } i = j \text{ then } x i \text{ else } 0)$

lemma diagvector-mult-vector[simp]: $\text{diagvector } x *v y = (\chi i. x i * y \$ i)$
unfolding diagvector-def matrix-vector-mult-def vec-eq-iff vec-lambda-beta
proof (rule, goal-cases)
case (1 i)
show ?case **by** (subst sum.remove[of - i], auto)

qed

lemma *diagvector-mult-left*: $\text{diagvector } x ** A = (\chi \ i \ j. x \ i * A \ \$ \ i \ \$ \ j)$ (**is** $?A = ?B$)

unfolding *vec-eq-iff*

proof (*intro allI*)

fix $i \ j$

show $?A \ \$ \ i \ \$ \ j = ?B \ \$ \ i \ \$ \ j$

unfolding *map-vector-def diagvector-def matrix-matrix-mult-def vec-lambda-beta*

by (*subst sum.remove[of - i], auto*)

qed

lemma *diagvector-mult-right*: $A ** \text{diagvector } x = (\chi \ i \ j. A \ \$ \ i \ \$ \ j * x \ j)$ (**is** $?A = ?B$)

unfolding *vec-eq-iff*

proof (*intro allI*)

fix $i \ j$

show $?A \ \$ \ i \ \$ \ j = ?B \ \$ \ i \ \$ \ j$

unfolding *map-vector-def diagvector-def matrix-matrix-mult-def vec-lambda-beta*

by (*subst sum.remove[of - j], auto*)

qed

lemma *diagvector-mult[simp]*: $\text{diagvector } x ** \text{diagvector } y = \text{diagvector } (\lambda \ i. x \ i * y \ i)$

unfolding *diagvector-mult-left* **unfolding** *diagvector-def* **by** (*auto simp: vec-eq-iff*)

lemma *diagvector-const[simp]*: $\text{diagvector } (\lambda \ x. k) = \text{mat } k$

unfolding *diagvector-def mat-def* **by** *auto*

lemma *diagvector-eq-mat*: $\text{diagvector } x = \text{mat } a \longleftrightarrow x = (\lambda \ x. a)$

unfolding *diagvector-def mat-def* **by** (*auto simp: vec-eq-iff*)

lemma *cmod-eq-Re*: **assumes** $\text{cmod } x = \text{Re } x$

shows *of-real* $(\text{Re } x) = x$

proof (*cases Im x = 0*)

case *False*

hence $(\text{cmod } x)^2 \neq (\text{Re } x)^2$ **unfolding** *norm-complex-def* **by** *simp*

from *this[unfolded assms]* **show** *?thesis* **by** *auto*

qed (*cases x, auto simp: norm-complex-def complex-of-real-def*)

hide-fact (**open**) *Matrix.vec-eq-iff*

no-notation

vec-index (**infixl** $\$ \ 100$)

lemma *spectral-radius-ev*:

$\exists \ ev \ v. \text{eigen-vector } A \ v \ ev \wedge \text{norm } ev = \text{spectral-radius } A$

proof –

from *non-empty-spectrum[of A]* *finite-spectrum[of A]* **have**

$spectral-radius\ A \in norm\ ' (Collect\ (eigen-value\ A))$
unfolding $spectral-radius-ev-def$ **by** $auto$
thus $?thesis\ unfolding\ eigen-value-def[abs-def]$ **by** $auto$
qed

lemma $spectral-radius-max$: **assumes** $eigen-value\ A\ v$
shows $norm\ v \leq spectral-radius\ A$
proof –
from $assms$ **have** $norm\ v \in norm\ ' (Collect\ (eigen-value\ A))$ **by** $auto$
from $Max-ge[OF\ -\ this,\ folded\ spectral-radius-ev-def]$
 $finite-spectrum[of\ A]$ **show** $?thesis$ **by** $auto$
qed

For Perron-Frobenius it is useful to use the linear norm, and not the Euclidean norm.

definition $norm1$:: $'a :: real-normed-field\ \hat{\ }^n \Rightarrow real$ **where**
 $norm1\ v = (\sum\ i \in UNIV.\ norm\ (v\ \$\ i))$

lemma $norm1-ge-0$: $norm1\ v \geq 0$ **unfolding** $norm1-def$
by $(rule\ sum-nonneg,\ auto)$

lemma $norm1-0[simp]$: $norm1\ 0 = 0$ **unfolding** $norm1-def$ **by** $auto$

lemma $norm1-nonzero$: **assumes** $v \neq 0$
shows $norm1\ v > 0$
proof –
from $\langle v \neq 0 \rangle$ **obtain** i **where** $vi: v\ \$\ i \neq 0$ **unfolding** $vec-eq-iff$
using $Finite-Cartesian-Product.vec-eq-iff\ zero-index$ **by** $force$
have $sum\ (\lambda\ i.\ norm\ (v\ \$\ i))\ (UNIV - \{i\}) \geq 0$
by $(rule\ sum-nonneg,\ auto)$
moreover **have** $norm\ (v\ \$\ i) > 0$ **using** vi **by** $auto$
ultimately
have $0 < norm\ (v\ \$\ i) + sum\ (\lambda\ i.\ norm\ (v\ \$\ i))\ (UNIV - \{i\})$ **by** $arith$
also **have** $\dots = norm1\ v$ **unfolding** $norm1-def$
by $(simp\ add:\ sum.remove)$
finally **show** $norm1\ v > 0$.
qed

lemma $norm1-0-iff[simp]$: $(norm1\ v = 0) = (v = 0)$
using $norm1-0\ norm1-nonzero$ **by** $(cases\ v = 0,\ force+)$

lemma $norm1-scaleR[simp]$: $norm1\ (r *_{R} v) = abs\ r * norm1\ v$ **unfolding** $norm1-def$
 $sum-distrib-left$
by $(rule\ sum.cong,\ auto)$

lemma $abs-norm1[simp]$: $abs\ (norm1\ v) = norm1\ v$ **using** $norm1-ge-0[of\ v]$ **by**
 $arith$

lemma $normalize-eigen-vector$: **assumes** $eigen-vector\ (A :: 'a :: real-normed-field)$

$\hat{\ }^n \hat{\ }^n) v \text{ ev}$
shows *eigen-vector* $A ((1 / \text{norm1 } v) *_R v) \text{ ev norm1 } ((1 / \text{norm1 } v) *_R v) = 1$
proof –
let $?v = (1 / \text{norm1 } v) *_R v$
from *assms*[*unfolded eigen-vector-def*]
have $\text{nz}: v \neq 0$ **and** $\text{id}: A *v v = \text{ev } *s v$ **by** *auto*
from nz **have** $\text{norm1}: \text{norm1 } v \neq 0$ **by** *auto*
thus $\text{norm1 } ?v = 1$ **by** *simp*
from $\text{norm1 } \text{nz}$ **have** $\text{nz}: ?v \neq 0$ **by** *auto*
have $A *v ?v = (1 / \text{norm1 } v) *_R (A *v v)$
by (*auto simp: vec-eq-iff matrix-vector-mult-def real-vector.scale-sum-right*)
also have $A *v v = \text{ev } *s v$ **unfolding** *id* ..
also have $(1 / \text{norm1 } v) *_R (\text{ev } *s v) = \text{ev } *s ?v$
by (*auto simp: vec-eq-iff*)
finally show *eigen-vector* $A ?v \text{ ev}$ **using** nz **unfolding** *eigen-vector-def* **by** *auto*
qed

lemma *norm1-cont*[*simp*]: *isCont* $\text{norm1 } v$ **unfolding** *norm1-def*[*abs-def*] **by** *auto*

lemma *norm1-ge-norm*: $\text{norm1 } v \geq \text{norm } v$ **unfolding** *norm1-def* *norm-vec-def*
by (*rule L2-set-le-sum, auto*)

The following continuity lemmas have been proven with hints from Fabian Immler.

lemma *tendsto-matrix-vector-mult*[*tendsto-intros*]:
 $((*v) (A :: 'a :: \text{real-normed-algebra-1 } \hat{\ }^n \hat{\ }^k) \longrightarrow A *v v)$ (*at* v *within* S)
unfolding *matrix-vector-mult-def*[*abs-def*]
by (*auto intro!: tendsto-intros*)

lemma *tendsto-matrix-matrix-mult*[*tendsto-intros*]:
 $((**) (A :: 'a :: \text{real-normed-algebra-1 } \hat{\ }^n \hat{\ }^k) \longrightarrow A ** B)$ (*at* B *within* S)
unfolding *matrix-matrix-mult-def*[*abs-def*]
by (*auto intro!: tendsto-intros*)

lemma *matrix-vec-scaleR*: $(A :: 'a :: \text{real-normed-algebra-1 } \hat{\ }^n \hat{\ }^k) *v (a *_R v)$
 $= a *_R (A *v v)$
unfolding *vec-eq-iff*
by (*auto simp: matrix-vector-mult-def scaleR-vec-def scaleR-sum-right*
intro!: sum.cong)

lemma (*in inj-semiring-hom*) *map-vector-0*: $(\text{map-vector } \text{hom } v = 0) = (v = 0)$
unfolding *vec-eq-iff* *map-vector-def* **by** *auto*

lemma (*in inj-semiring-hom*) *map-vector-inj*: $(\text{map-vector } \text{hom } v = \text{map-vector } \text{hom } w) = (v = w)$
unfolding *vec-eq-iff* *map-vector-def* **by** *auto*

lemma (*in semiring-hom*) *matrix-vector-mult-hom*:

$(\text{map-matrix hom } A) * v (\text{map-vector hom } v) = \text{map-vector hom } (A * v v)$
by (transfer fixing: hom, auto simp: mult-mat-vec-hom)

lemma (in semiring-hom) vector-smult-hom:
 $\text{hom } x * s (\text{map-vector hom } v) = \text{map-vector hom } (x * s v)$
by (transfer fixing: hom, auto simp: vec-hom-smult)

lemma (in inj-comm-ring-hom) eigen-vector-hom:
 $\text{eigen-vector } (\text{map-matrix hom } A) (\text{map-vector hom } v) (\text{hom } x) = \text{eigen-vector } A$
 $v x$
unfolding eigen-vector-def matrix-vector-mult-hom vector-smult-hom map-vector-0
map-vector-inj
by auto

end

4.2 Perron-Frobenius theorem via Brouwer's fixpoint theorem.

theory Perron-Frobenius
imports
HOL-Analysis.Brouwer-Fixpoint
Perron-Frobenius-Aux
begin

We follow the textbook proof of Serre [2, Theorem 5.2.1].

context
fixes $A :: \text{complex } ^\wedge n ^\wedge n :: \text{finite}$
assumes $\text{rnn}A: \text{real-non-neg-mat } A$
begin

private abbreviation (input) sr **where** $sr \equiv \text{spectral-radius } A$

private definition $\text{max-v-ev} :: (\text{complex } ^\wedge n) \times \text{complex}$ **where**
 $\text{max-v-ev} = (\text{SOME } v\text{-ev. eigen-vector } A (\text{fst } v\text{-ev}) (\text{snd } v\text{-ev}))$
 $\wedge \text{norm } (\text{snd } v\text{-ev}) = sr$

private definition $\text{max-v} = (1 / \text{norm1 } (\text{fst } \text{max-v-ev})) *_R \text{fst } \text{max-v-ev}$

private definition $\text{max-ev} = \text{snd } \text{max-v-ev}$

private lemma max-v-ev :
 $\text{eigen-vector } A \text{max-v } \text{max-ev}$
 $\text{norm } \text{max-ev} = sr$
 $\text{norm1 } \text{max-v} = 1$

proof –

obtain $v \text{ ev}$ **where** $\text{id}: \text{max-v-ev} = (v, \text{ev})$ **by** force
from $\text{spectral-radius-ev}[\text{of } A]$ $\text{someI-ex}[\text{of } \lambda v\text{-ev. eigen-vector } A (\text{fst } v\text{-ev}) (\text{snd } v\text{-ev})]$
 $\wedge \text{norm } (\text{snd } v\text{-ev}) = sr, \text{folded } \text{max-v-ev-def}, \text{unfolded id}]$

```

have  $v$ : eigen-vector  $A$   $v$   $ev$  and  $ev$ : norm  $ev = sr$  by auto
from normalize-eigen-vector[ $OF$   $v$ ]  $ev$ 
show eigen-vector  $A$   $max\text{-}v$   $max\text{-}ev$  norm  $max\text{-}ev = sr$  norm1  $max\text{-}v = 1$ 
  unfolding max-v-def max-ev-def id by auto
qed

```

In the definition of S , we use the linear norm instead of the default euclidean norm which is defined via the type-class. The reason is that S is not convex if one uses the euclidean norm.

```

private definition  $B$  ::  $real$   $^n$   $^n$  where  $B \equiv \chi$   $i$   $j$ . Re ( $A$   $\$$   $i$   $\$$   $j$ )
private definition  $S$  where  $S = \{v :: real$   $^n$  . norm1  $v = 1 \wedge (\forall i. v$   $\$$   $i \geq 0) \wedge$ 
  ( $\forall i. (B * v$   $v) \$ i \geq sr * (v$   $\$ i))\}$ 
private definition  $f$  ::  $real$   $^n \Rightarrow real$   $^n$  where
   $f v = (1 / norm1 (B * v v)) *_R (B * v v)$ 

```

```

private lemma closedS: closed  $S$ 
  unfolding S-def matrix-vector-mult-def[abs-def]
proof (intro closed-Collect-conj closed-Collect-all closed-Collect-le closed-Collect-eq)
  show continuous-on  $UNIV$  norm1
  by (simp add: continuous-at-imp-continuous-on)
qed (auto intro!: continuous-intros continuous-on-component)

```

```

private lemma boundedS: bounded  $S$ 
proof -
  {
    fix  $v :: real$   $^n$ 
    from norm1-ge-norm[of  $v$ ] have norm1  $v = 1 \implies norm$   $v \leq 1$  by auto
  }
  thus ?thesis
  unfolding S-def bounded-iff
  by (auto intro!: exI[of - 1])
qed

```

```

private lemma compactS: compact  $S$ 
  using boundedS closedS
  by (simp add: compact-eq-bounded-closed)

```

```

private lemmas rnn = real-non-neg-matD[ $OF$  rnnA]

```

```

lemma B-norm:  $B$   $\$$   $i$   $\$$   $j = norm$  ( $A$   $\$$   $i$   $\$$   $j$ )
  using rnn[of  $i$   $j$ ]
  by (cases  $A$   $\$$   $i$   $\$$   $j$ , auto simp: B-def)

```

```

lemma mult-B-mono: assumes  $\bigwedge i. v$   $\$$   $i \geq w$   $\$$   $i$ 
  shows  $(B * v v) \$ i \geq (B * w w) \$ i$  unfolding matrix-vector-mult-def vec-lambda-beta
  by (rule sum-mono, rule mult-left-mono[ $OF$  assms], unfold B-norm, auto)

```

```

private lemma non-emptyS: S ≠ {}
proof -
  let ?v = (χ i. norm (max-v $ i)) :: real ^ 'n
  have norm1 max-v = 1 by (rule max-v-ev(3))
  hence nv: norm1 ?v = 1 unfolding norm1-def by auto
  {
    fix i
    have sr * (?v $ i) = sr * norm (max-v $ i) by auto
    also have ... = (norm max-v) * norm (max-v $ i) using max-v-ev by auto
    also have ... = norm ((max-v *s max-v) $ i) by (auto simp: norm-mult)
    also have max-v *s max-v = A *v max-v using max-v-ev(1)[unfolded eigen-vector-def]
  }
by auto
  also have norm ((A *v max-v) $ i) ≤ (B *v ?v) $ i
  unfolding matrix-vector-mult-def vec-lambda-beta
  by (rule sum-norm-le, auto simp: norm-mult B-norm)
  finally have sr * (?v $ i) ≤ (B *v ?v) $ i .
} note le = this
have ?v ∈ S unfolding S-def using nv le by auto
thus ?thesis by blast
qed

```

```

private lemma convexS: convex S
proof (rule convexI)
  fix v w a b
  assume *: v ∈ S w ∈ S 0 ≤ a 0 ≤ b a + b = (1 :: real)
  let ?lin = a *R v + b *R w
  from * have 1: norm1 v = 1 norm1 w = 1 unfolding S-def by auto
  have norm1 ?lin = a * norm1 v + b * norm1 w
  unfolding norm1-def sum-distrib-left sum.distrib[symmetric]
proof (rule sum.cong)
  fix i :: 'n
  from * have v $ i ≥ 0 w $ i ≥ 0 unfolding S-def by auto
  thus norm (?lin $ i) = a * norm (v $ i) + b * norm (w $ i)
  using *(3-4) by auto
qed simp
also have ... = 1 using *(5) 1 by auto
finally have norm1: norm1 ?lin = 1 .
{
  fix i
  from * have 0 ≤ v $ i sr * v $ i ≤ (B *v v) $ i unfolding S-def by auto
  with ⟨a ≥ 0⟩ have a: a * (sr * v $ i) ≤ a * (B *v v) $ i by (intro mult-left-mono)
  from * have 0 ≤ w $ i sr * w $ i ≤ (B *v w) $ i unfolding S-def by auto
  with ⟨b ≥ 0⟩ have b: b * (sr * w $ i) ≤ b * (B *v w) $ i by (intro
mult-left-mono)
  from a b have a * (sr * v $ i) + b * (sr * w $ i) ≤ a * (B *v v) $ i + b *
(B *v w) $ i by auto
} note le = this
have switch[simp]: ∧ x y. x * a * y = a * x * y ∧ x y. x * b * y = b * x * y
by auto

```

have [simp]: $x \in \{v, w\} \implies a * (r * x \$ h i) = r * (a * x \$ h i)$ **for** $a r i x$ **by**
auto
show $a *_R v + b *_R w \in S$ **using** * norm1 le **unfolding** S-def
by (*auto simp: matrix-vect-scaleR matrix-vector-right-distrib ring-distrib*)
qed

private abbreviation (input) $r :: \text{real} \Rightarrow \text{complex}$ **where**
 $r \equiv \text{of-real}$

private abbreviation $rv :: \text{real} \wedge^n \Rightarrow \text{complex} \wedge^n$ **where**
 $rv v \equiv \chi i. r (v \$ i)$

private lemma $rv-0$: $(rv v = 0) = (v = 0)$
by (*simp add: of-real-hom.map-vector-0 map-vector-def vec-eq-iff*)

private lemma $rv-mult$: $A * v rv v = rv (B * v v)$
proof –
have *map-matrix* $r B = A$
using *rnnA* **unfolding** *map-matrix-def B-def real-non-neg-mat-def map-vector-def elements-mat-h-def*
by *vector*
thus *?thesis*
using *of-real-hom.matrix-vector-mult-hom*[*of B, where 'a = complex*]
unfolding *map-vector-def* **by** *auto*
qed

context
assumes $zero-no-ev$: $\bigwedge v. v \in S \implies A * v rv v \neq 0$
begin
private lemma $normB-S$: **assumes** $v: v \in S$
shows $norm1 (B * v v) \neq 0$
proof –
from $zero-no-ev$ [*OF v, unfolded rv-mult rv-0*]
show *?thesis* **by** *auto*
qed

private lemma $image-f$: $f ' S \subseteq S$
proof –
{
fix v
assume $v: v \in S$
hence $norm: norm1 v = 1$ **and** $ge: \bigwedge i. v \$ i \geq 0 \wedge i. sr * v \$ i \leq (B * v v)$
 $\$ i$ **unfolding** S-def **by** *auto*
from $normB-S$ [*OF v*] **have** $normB: norm1 (B * v v) > 0$ **using** $norm1-nonzero$
by *auto*
have $fv: f v = (1 / norm1 (B * v v)) *_R (B * v v)$ **unfolding** f-def **by** *auto*
from $normB$ **have** $Bv0: B * v v \neq 0$ **unfolding** $norm1-0-iff$ [*symmetric*] **by**
linarith
have $norm: norm1 (f v) = 1$ **unfolding** fv **using** $normB Bv0$ **by** *simp*

```

define c where c = (1 / norm1 (B *v v))
have c: c > 0 unfolding c-def using normB by auto
{
  fix i
  have 1: f v $ i ≥ 0 unfolding fv c-def[symmetric] using c ge
  by (auto simp: matrix-vector-mult-def sum-distrib-left B-norm intro!: sum-nonneg)
  have id1:  $\bigwedge i. (B *v f v) $ i = c * ((B *v (B *v v)) $ i)$ 
    unfolding f-def c-def matrix-vect-scaleR by simp
  have id3:  $\bigwedge i. sr * f v $ i = c * ((B *v (sr *_R v)) $ i)$ 
    unfolding f-def c-def[symmetric] matrix-vect-scaleR by auto
  have 2:  $sr * f v $ i \leq (B *v f v) $ i$  unfolding id1 id3
    unfolding mult-le-cancel-iff2[OF ⟨c > 0⟩]
    by (rule mult-B-mono, insert ge(2), auto)
  note 1 2
}
with norm have f v ∈ S unfolding S-def by auto
}
thus ?thesis by blast
qed

```

```

private lemma cont-f: continuous-on S f
  unfolding f-def[abs-def] continuous-on using normB-S
  unfolding norm1-def
  by (auto intro!: tendsto-eq-intros)

```

qualified lemma perron-frobenius-positive-ev:

$\exists v. \text{eigen-vector } A \ v \ (r \ sr) \wedge \text{real-non-neg-vec } v$

proof –

```

from brouwer[OF compactS convexS non-emptyS cont-f image-f]
  obtain v where v: v ∈ S and fv: f v = v by auto
define ev where ev = norm1 (B *v v)
from normB-S[OF v] have ev ≠ 0 unfolding ev-def by auto
with norm1-ge-0[of B *v v, folded ev-def] have norm: ev > 0 by auto
from arg-cong[OF fv[unfolded f-def], of λ (w :: real ^ 'n). ev *_R w] norm
  have ev: B *v v = ev *s v unfolding ev-def[symmetric] scalar-mult-eq-scaleR
by simp
with v[unfolded S-def] have ge:  $\bigwedge i. sr * v $ i \leq ev * v $ i$  by auto
have A *v rv v = rv (B *v v) unfolding rv-mult ..
also have ... = ev *s rv v unfolding ev vec-eq-iff
  by (simp add: scaleR-conv-of-real scaleR-vec-def)
finally have ev: A *v rv v = ev *s rv v .
from v have v0: v ≠ 0 unfolding S-def by auto
hence rv v ≠ 0 unfolding rv-0 .
with ev have ev: eigen-vector A (rv v) ev unfolding eigen-vector-def by auto
hence eigen-value A ev unfolding eigen-value-def by auto
from spectral-radius-max[OF this] have le: norm (r ev) ≤ sr .
from v0 obtain i where v $ i ≠ 0 unfolding vec-eq-iff by auto
from v have v $ i ≥ 0 unfolding S-def by auto
with ⟨v $ i ≠ 0⟩ have v $ i > 0 by auto

```

```

with  $ge[of\ i]$  have  $ge: sr \leq ev$  by auto
with  $le$  have  $sr: r\ sr = ev$  by auto
from  $v$  have  $*$ : real-non-neg-vec  $(rv\ v)$  unfolding  $S-def$  real-non-neg-vec-def
vec-elements-h-def by auto
show ?thesis unfolding  $sr$ 
by (rule  $exI[of\ -\ rv\ v]$ , insert  $*\ ev\ norm$ , auto)
qed
end

```

```

qualified lemma perron-frobenius-both:
 $\exists v.$  eigen-vector  $A\ v\ (r\ sr) \wedge$  real-non-neg-vec  $v$ 
proof (cases  $\forall v \in S. A *v\ rv\ v \neq 0$ )
case True
show ?thesis
by (rule Perron-Frobenius.perron-frobenius-positive-ev[OF rnnA], insert True,
auto)
next
case False
then obtain  $v$  where  $v: v \in S$  and  $A0: A *v\ rv\ v = 0$  by auto
hence  $id: A *v\ rv\ v = 0 *s\ rv\ v$  and  $v0: v \neq 0$  unfolding  $S-def$  by auto
from  $v0$  have  $rv\ v \neq 0$  unfolding  $rv-0$  .
with  $id$  have  $ev: eigen-vector\ A\ (rv\ v)\ 0$  unfolding eigen-vector-def by auto
hence eigen-value  $A\ 0$  unfolding eigen-value-def ..
from spectral-radius-max[OF this] have  $0: 0 \leq sr$  by auto
from  $v[unfolding\ S-def]$  have  $ge: \bigwedge i. sr * v\ \$\ i \leq (B *v\ v)\ \$\ i$  by auto
from  $v[unfolding\ S-def]$  have  $rnn: real-non-neg-vec\ (rv\ v)$ 
unfolding real-non-neg-vec-def vec-elements-h-def by auto
from  $v0$  obtain  $i$  where  $v\ \$\ i \neq 0$  unfolding vec-eq-iff by auto
from  $v$  have  $v\ \$\ i \geq 0$  unfolding  $S-def$  by auto
with  $\langle v\ \$\ i \neq 0 \rangle$  have  $vi: v\ \$\ i > 0$  by auto
from  $rv-mult[of\ v, unfolded\ A0]$  have  $rv\ (B *v\ v) = 0$  by simp
hence  $B *v\ v = 0$  unfolding  $rv-0$  .
from  $ge[of\ i, unfolded\ this]\ vi$  have  $ge: sr \leq 0$  by (simp add: mult-le-0-iff)
with  $\langle 0 \leq sr \rangle$  have  $sr = 0$  by auto
show ?thesis unfolding  $\langle sr = 0 \rangle$  using  $rnn\ ev$  by auto
qed
end

```

Perron Frobenius: The largest complex eigenvalue of a real-valued non-negative matrix is a real one, and it has a real-valued non-negative eigenvector.

```

lemma perron-frobenius:
assumes real-non-neg-mat  $A$ 
shows  $\exists v.$  eigen-vector  $A\ v\ (of-real\ (spectral-radius\ A)) \wedge$  real-non-neg-vec  $v$ 
by (rule Perron-Frobenius.perron-frobenius-both[OF assms])

```

And a version which ignores the eigenvector.

```

lemma perron-frobenius-eigen-value:
assumes real-non-neg-mat  $A$ 

```

shows eigen-value A (of-real (spectral-radius A))
 using perron-frobenius[OF assms] unfolding eigen-value-def by blast

end

5 Roots of Unity

theory Roots-Unity

imports

Polynomial-Factorization.Order-Polynomial
 HOL-Computational-Algebra.Fundamental-Theorem-Algebra
 Polynomial-Interpolation.Ring-Hom-Poly

begin

lemma cis-mult-cmod-id: $\text{cis } (Arg\ x) * \text{of-real } (cmod\ x) = x$
 using cis-cmod-Arg[unfolded rcis-def] by (simp add: ac-simps)

lemma rcis-mult-cis[simp]: $\text{rcis } n\ a * \text{cis } b = \text{rcis } n\ (a + b)$ unfolding cis-rcis-eq
 rcis-mult by simp

lemma rcis-div-cis[simp]: $\text{rcis } n\ a / \text{cis } b = \text{rcis } n\ (a - b)$ unfolding cis-rcis-eq
 rcis-divide by simp

lemma cis-plus-2pi[simp]: $\text{cis } (x + 2 * pi) = \text{cis } x$ by (auto simp: complex-eq-iff)

lemma cis-plus-2pi-neq-1: assumes $x: 0 < x < 2 * pi$

shows $\text{cis } x \neq 1$

proof -

from x have $\cos\ x \neq 1$ by (smt cos-2pi-minus cos-monotone-0-pi cos-zero)

thus ?thesis by (auto simp: complex-eq-iff)

qed

lemma cis-times-2pi[simp]: $\text{cis } (of\text{-nat } n * 2 * pi) = 1$

proof (induct n)

case (Suc n)

have $of\text{-nat } (Suc\ n) * 2 * pi = of\text{-nat } n * 2 * pi + 2 * pi$ by (simp add:
 distrib-right)

also have $\text{cis } \dots = 1$ unfolding cis-plus-2pi Suc ..

finally show ?case .

qed simp

lemma cis-add-pi[simp]: $\text{cis } (pi + x) = - \text{cis } x$

by (auto simp: complex-eq-iff)

lemma cis-3-pi-2[simp]: $\text{cis } (pi * 3 / 2) = - i$

proof -

have $\text{cis } (pi * 3 / 2) = \text{cis } (pi + pi / 2)$

by (rule arg-cong[of - - cis], simp)

also have $\dots = - i$ unfolding cis-add-pi by simp

finally show ?thesis .

qed

lemma *rcis-plus-2pi*[simp]: $rcis\ y\ (x + 2 * pi) = rcis\ y\ x$ **unfolding** *rcis-def* **by** *simp*

lemma *rcis-times-2pi*[simp]: $rcis\ r\ (of\ nat\ n * 2 * pi) = of\ real\ r$
unfolding *rcis-def cis-times-2pi* **by** *simp*

lemma *arg-rcis-cis*: **assumes** $n: n > 0$ **shows** $Arg\ (rcis\ n\ x) = Arg\ (cis\ x)$
using *Arg-bounded cis-Arg-unique cis-Arg complex-mod-rcis n rcis-def sgn-eq* **by** *auto*

lemma *arg-eqD*: **assumes** $Arg\ (cis\ x) = Arg\ (cis\ y) -pi < x \leq pi -pi < y \leq pi$
shows $x = y$
using *assms(1)* **unfolding** *cis-Arg-unique[OF sgn-cis assms(2-3)] cis-Arg-unique[OF sgn-cis assms(4-5)]* .

lemma *rcis-inj-on*: **assumes** $r: r \neq 0$ **shows** *inj-on* $(rcis\ r)\ \{0 ..< 2 * pi\}$
proof (*rule inj-onI, goal-cases*)
case $(1\ x\ y)$
from *arg-cong[OF 1(3), of $\lambda x. x / r$]* **have** $cis\ x = cis\ y$ **using** r **by** (*simp add: rcis-def*)
from *arg-cong[OF this, of $\lambda x. inverse\ x$]* **have** $cis\ (-x) = cis\ (-y)$ **by** *simp*
from *arg-cong[OF this, of uminus]* **have** $*$: $cis\ (-x + pi) = cis\ (-y + pi)$
by (*auto simp: complex-eq-iff*)
have $-x + pi = -y + pi$
by (*rule arg-eqD[OF arg-cong[OF *, of Arg]]*, *insert 1(1-2)*, *auto*)
thus $?case$ **by** *simp*
qed

lemma *cis-inj-on*: *inj-on* $cis\ \{0 ..< 2 * pi\}$
using *rcis-inj-on[of 1]* **unfolding** *rcis-def* **by** *auto*

definition *root-unity* :: $nat \Rightarrow 'a :: comm-ring-1\ poly$ **where**
root-unity $n = monom\ 1\ n - 1$

lemma *poly-root-unity*: *poly* $(root-unity\ n)\ x = 0 \longleftrightarrow x^n = 1$
unfolding *root-unity-def* **by** (*simp add: poly-monom*)

lemma *degree-root-unity*[simp]: *degree* $(root-unity\ n) = n$ (**is** *degree* $?p = -$)
proof –
have $p: ?p = monom\ 1\ n + (-1)$ **unfolding** *root-unity-def* **by** *auto*
show $?thesis$
proof (*cases n*)
case 0
thus $?thesis$ **unfolding** p **by** *simp*
next
case $(Suc\ m)$
show $?thesis$ **unfolding** p **unfolding** *Suc*
by (*subst degree-add-eq-left, auto simp: degree-monom-eq*)

qed
qed

lemma *zero-root-unity*[simp]: $\text{root-unity } n = 0 \iff n = 0$ (is $?p = 0 \iff -$)

proof (cases $n = 0$)

case *True*

thus *?thesis* unfolding *root-unity-def* by *simp*

next

case *False*

from *degree-root-unity*[of n] *False*

have $\text{degree } ?p \neq 0$ by *auto*

hence $?p \neq 0$ by *fastforce*

thus *?thesis* using *False* by *auto*

qed

definition *prod-root-unity* :: $\text{nat list} \Rightarrow 'a :: \text{idom poly}$ where

prod-root-unity ns = prod-list (map root-unity ns)

lemma *poly-prod-root-unity*: $\text{poly } (\text{prod-root-unity } ns) x = 0 \iff (\exists k \in \text{set } ns. x^{\wedge k} = 1)$

unfolding *prod-root-unity-def*

by (*simp add: poly-prod-list prod-list-zero-iff o-def image-def poly-root-unity*)

lemma *degree-prod-root-unity*[simp]: $0 \notin \text{set } ns \implies \text{degree } (\text{prod-root-unity } ns) = \text{sum-list } ns$

unfolding *prod-root-unity-def*

by (*subst degree-prod-list-eq, auto simp: o-def*)

lemma *zero-prod-root-unity*[simp]: $\text{prod-root-unity } ns = 0 \iff 0 \in \text{set } ns$

unfolding *prod-root-unity-def prod-list-zero-iff* by *auto*

lemma *roots-of-unity*: **assumes** $n: n \neq 0$

shows $(\lambda i. (\text{cis } (\text{of-nat } i * 2 * \text{pi} / n))) \text{ ' } \{0 ..< n\} = \{x :: \text{complex}. x^{\wedge n} = 1\}$ (is *?prod = ?Roots*)

$\{x. \text{poly } (\text{root-unity } n) x = 0\} = \{x :: \text{complex}. x^{\wedge n} = 1\}$

$\text{card } \{x :: \text{complex}. x^{\wedge n} = 1\} = n$

proof (*atomize(full), goal-cases*)

case 1

let $?one = 1 :: \text{complex}$

let $?p = \text{monom } ?one n - 1$

have *degM*: $\text{degree } (\text{monom } ?one n) = n$ by (*rule degree-monom-eq, simp*)

have $\text{degree } ?p = \text{degree } (\text{monom } ?one n + (-1))$ by *simp*

also have $\dots = \text{degree } (\text{monom } ?one n)$

by (*rule degree-add-eq-left, insert n, simp add: degM*)

finally have *degp*: $\text{degree } ?p = n$ unfolding *degM* .

with n have $p: ?p \neq 0$ by *auto*

have *roots*: $?Roots = \{x. \text{poly } ?p x = 0\}$

unfolding *poly-diff poly-monom* by *simp*

also have *finite* \dots by (*rule poly-roots-finite[OF p]*)

```

finally have fin: finite ?Roots .
have sub: ?prod  $\subseteq$  ?Roots
proof
  fix x
  assume  $x \in ?prod$ 
  then obtain i where  $x = cis (real\ i * 2 * pi / n)$  by auto
  have  $x^n = cis (real\ i * 2 * pi)$  unfolding x DeMoirve using n by simp
  also have  $\dots = 1$  by simp
  finally show  $x \in ?Roots$  by auto
qed
have Rn:  $card\ ?Roots \leq n$  unfolding roots
  by (rule poly-roots-degree[of ?p, unfolded degp, OF p])
have  $\dots = card\ \{0 ..< n\}$  by simp
also have  $\dots = card\ ?prod$ 
proof (rule card-image[symmetric], rule inj-onI, goal-cases)
  case (1 x y)
  {
    fix m
    assume  $m < n$ 
    hence  $real\ m < real\ n$  by simp
    from mult-strict-right-mono[OF this, of 2 * pi / real n] n
    have  $real\ m * 2 * pi / real\ n < real\ n * 2 * pi / real\ n$  by simp
    hence  $real\ m * 2 * pi / real\ n < 2 * pi$  using n by simp
  } note [simp] = this
  have  $0: (1 :: real) \neq 0$  using n by auto
  have  $real\ x * 2 * pi / real\ n = real\ y * 2 * pi / real\ n$ 
    by (rule inj-onD[OF rcis-inj-on 1(3)[unfolded cis-rcis-eq]], insert 1(1-2),
auto)
  with n show  $x = y$  by auto
qed
finally have cn:  $card\ ?prod = n$  ..
with Rn have  $card\ ?prod \geq card\ ?Roots$  by auto
with card-mono[OF fin sub] have card:  $card\ ?prod = card\ ?Roots$  by auto
have  $?prod = ?Roots$ 
  by (rule card-subset-eq[OF fin sub card])
  from this roots[symmetric] cn[unfolded this]
  show ?case unfolding root-unity-def by blast
qed

lemma poly-roots-dvd: fixes  $p :: 'a :: field\ poly$ 
  assumes  $p \neq 0$  and  $degree\ p = n$ 
  and  $card\ \{x. poly\ p\ x = 0\} \geq n$  and  $\{x. poly\ p\ x = 0\} \subseteq \{x. poly\ q\ x = 0\}$ 
shows  $p\ dvd\ q$ 
proof -
  from poly-roots-degree[OF assms(1)] assms(2-3) have  $card\ \{x. poly\ p\ x = 0\}$ 
  = n by auto
  from assms(1-2) this assms(4)
  show ?thesis
  proof (induct n arbitrary: p q)

```

```

    case (0 p q)
    from is-unit-iff-degree[OF 0(1)] 0(2) show ?case by blast
next
case (Suc n p q)
let ?P = {x. poly p x = 0}
let ?Q = {x. poly q x = 0}
from Suc(4-5) card-gt-0-iff[of ?P] obtain x where
  x: poly p x = 0 poly q x = 0 and fin: finite ?P by auto
define r where r = [:-x, 1:]
from x[unfolded poly-eq-0-iff-dvd r-def[symmetric]] obtain p' q' where
  p: p = r * p' and q: q = r * q' unfolding dvd-def by auto
from Suc(2) have degree p = degree r + degree p' unfolding p
  by (subst degree-mult-eq, auto)
with Suc(3) have deg: degree p' = n unfolding r-def by auto
from Suc(2) p have p'0: p' ≠ 0 by auto
let ?P' = {x. poly p' x = 0}
let ?Q' = {x. poly q' x = 0}
have P: ?P = insert x ?P' unfolding p poly-mult unfolding r-def by auto
have Q: ?Q = insert x ?Q' unfolding q poly-mult unfolding r-def by auto
{
  assume x ∈ ?P'
  hence ?P = ?P' unfolding P by auto
  from arg-cong[OF this, of card, unfolded Suc(4)] deg have False
    using poly-roots-degree[OF p'0] by auto
} note xp' = this
hence xP': x ∉ ?P' by auto
have card ?P = Suc (card ?P') unfolding P
  by (rule card-insert-disjoint[OF - xP'], insert fin[unfolded P], auto)
with Suc(4) have card: card ?P' = n by auto
from Suc(5)[unfolded P Q] xP' have ?P' ⊆ ?Q' by auto
from Suc(1)[OF p'0 deg card this]
have IH: p' dvd q' .
show ?case unfolding p q using IH by simp
qed
qed

```

lemma *root-unity-decomp*: assumes $n: n \neq 0$

shows *root-unity* $n =$

prod-list (map ($\lambda i. [:-cis (of-nat i * 2 * pi / n), 1:] [0 ..< n]$) (is ?u = ?p))

proof –

have *deg*: degree ?u = n by simp

note *main* = *roots-of-unity*[OF n]

have *dvd*: ?u dvd ?p

proof (rule *poly-roots-dvd*[OF - *deg*])

show $n \leq \text{card } \{x. \text{poly } ?u x = 0\}$ using *main* by auto

show ?u ≠ 0 using n by auto

show $\{x. \text{poly } ?u x = 0\} \subseteq \{x. \text{poly } ?p x = 0\}$

unfolding *main*(2) *main*(1)[*symmetric*] *poly-prod-list prod-list-zero-iff* by

auto

```

qed
have deg': degree ?p = n
  by (subst degree-prod-list-eq, auto simp: o-def sum-list-triv)
have mon: monic ?u using deg unfolding root-unity-def using n by auto
have mon': monic ?p by (rule monic-prod-list, auto)
from dvd[unfolded dvd-def] obtain f where puf: ?p = ?u * f by auto
have degree ?p = degree ?u + degree f using mon' n unfolding puf
  by (subst degree-mult-eq, auto)
with deg deg' have degree f = 0 by auto
from degree0-coeffs[OF this] obtain a where f: f = [:a:] by blast
from arg-cong[OF puf, of lead-coeff] mon mon'
have a = 1 unfolding puf f by (cases a = 0, auto)
with f have f: f = 1 by auto
with puf show ?thesis by auto
qed

lemma order-monic-linear: order x [:y,1:] = (if y + x = 0 then 1 else 0)
proof (cases y + x = 0)
  case True
  hence poly [:y,1:] x = 0 by simp
  from this[unfolded order-root] have order x [:y,1:] ≠ 0 by auto
  moreover from order-degree[of [:y,1:] x] have order x [:y,1:] ≤ 1 by auto
  ultimately show ?thesis unfolding True by auto
next
  case False
  hence poly [:y,1:] x ≠ 0 by auto
  from order-0I[OF this] False show ?thesis by auto
qed

lemma order-root-unity: fixes x :: complex assumes n: n ≠ 0
  shows order x (root-unity n) = (if x^n = 1 then 1 else 0)
  (is order - ?u = -)
proof (cases x^n = 1)
  case False
  with roots-of-unity(2)[OF n] have poly ?u x ≠ 0 by auto
  from False order-0I[OF this] show ?thesis by auto
next
  case True
  let ?phi = λ i :: nat. i * 2 * pi / n
  from True roots-of-unity(1)[OF n] obtain i where i: i < n
    and x: x = cis (?phi i) by force
  from i have n-split: [0 ..< n] = [0 ..< i] @ i # [Suc i ..< n]
    by (metis le-Suc-ex less-imp-le-nat not-le-imp-less not-less0 upt-add-eq-append
    upt-conv-Cons)
  {
    fix j
    assume j: j < n ∨ j < i and eq: cis (?phi i) = cis (?phi j)
    from inj-onD[OF cis-inj-on eq] i j n have i = j by (auto simp: field-simps)
  } note inj = this

```

have $order\ x\ ?u = 1$ **unfolding** *root-unity-decomp*[*OF n*]
unfolding $x\ n\text{-split}$ **using** *inj*
by (*subst order-prod-list, force, fastforce simp: order-monic-linear*)
with *True* **show** *?thesis* **by** *auto*
qed

lemma *order-prod-root-unity*: **assumes** $0 \notin set\ ks$
shows $order\ (x :: complex)\ (prod\text{-root-unity}\ ks) = length\ (filter\ (\lambda\ k.\ x^k = 1)\ ks)$
proof –
have $order\ x\ (prod\text{-root-unity}\ ks) = (\sum\ k \leftarrow ks.\ order\ x\ (root\text{-unity}\ k))$
unfolding *prod-root-unity-def*
by (*subst order-prod-list, insert 0, auto simp: o-def*)
also have $\dots = (\sum\ k \leftarrow ks.\ (if\ x^k = 1\ then\ 1\ else\ 0))$
by (*rule arg-cong, rule map-cong, insert 0, force, intro order-root-unity, metis*)
also have $\dots = length\ (filter\ (\lambda\ k.\ x^k = 1)\ ks)$
by (*subst sum-list-map-filter^[symmetric], simp add: sum-list-triv*)
finally show *?thesis* .
qed

lemma *root-unity-witness*: **fixes** $xs :: complex\ list$
assumes $prod\text{-list}\ (map\ (\lambda\ x.\ [-x, 1:])\ xs) = monom\ 1\ n - 1$
shows $x^n = 1 \iff x \in set\ xs$
proof –
from *assms* **have** $n \neq 0$ **by** (*cases n = 0, auto simp: prod-list-zero-iff*)
have $x \in set\ xs \iff poly\ (prod\text{-list}\ (map\ (\lambda\ x.\ [-x, 1:])\ xs))\ x = 0$
unfolding *poly-prod-list prod-list-zero-iff* **by** *auto*
also have $\dots \iff x^n = 1$ **using** *roots-of-unity(2)*[*OF n0*] **unfolding** *assms*
root-unity-def **by** *auto*
finally show *?thesis* **by** *auto*
qed

lemma *root-unity-explicit*: **fixes** $x :: complex$
shows
 $(x^1 = 1) \iff x = 1$
 $(x^2 = 1) \iff (x \in \{1, -1\})$
 $(x^3 = 1) \iff (x \in \{1, Complex\ (-1/2)\ (sqrt\ 3 / 2), Complex\ (-1/2)\ (-sqrt\ 3 / 2)\})$
 $(x^4 = 1) \iff (x \in \{1, -1, i, -i\})$
proof –
show $(x^1 = 1) \iff x = 1$
by (*subst root-unity-witness*[*of [1]*], *code-simp, auto*)
show $(x^2 = 1) \iff (x \in \{1, -1\})$
by (*subst root-unity-witness*[*of [1, -1]*], *code-simp, auto*)
show $(x^4 = 1) \iff (x \in \{1, -1, i, -i\})$
by (*subst root-unity-witness*[*of [1, -1, i, -i]*], *code-simp, auto*)
have $3 = Suc\ (Suc\ (Suc\ 0))\ 1 = [1:]$ **by** *auto*
show $(x^3 = 1) \iff (x \in \{1, Complex\ (-1/2)\ (sqrt\ 3 / 2), Complex\ (-1/2)\ (-sqrt\ 3 / 2)\})$

by (subst root-unity-witness[of
 [1, Complex (-1/2) (sqrt 3 / 2), Complex (-1/2) (- sqrt 3 / 2)]],
 auto simp: 3 monom-altdef complex-mult complex-eq-iff)
 qed

definition primitive-root-unity :: nat \Rightarrow 'a :: power \Rightarrow bool **where**
 primitive-root-unity k x = (k \neq 0 \wedge x^k = 1 \wedge (\forall k' < k. k' \neq 0 \longrightarrow x^{k'} \neq 1))

lemma primitive-root-unityD: **assumes** primitive-root-unity k x
 shows k \neq 0 x^k = 1 k' \neq 0 \implies x^{k'} = 1 \implies k \leq k'
proof -
 note * = assms[unfolded primitive-root-unity-def]
 from * have **: k' < k \implies k' \neq 0 \implies x^{k'} \neq 1 by auto
 show k \neq 0 x^k = 1 using * by auto
 show k' \neq 0 \implies x^{k'} = 1 \implies k \leq k' using ** by force
 qed

lemma primitive-root-unity-exists: **assumes** k \neq 0 x^k = 1
 shows \exists k'. k' \leq k \wedge primitive-root-unity k' x
proof -
 let ?P = λ k. x^k = 1 \wedge k \neq 0
 define k' **where** k' = (LEAST k. ?P k)
 from assms **have** Pk: \exists k. ?P k by auto
 from LeastI-ex[OF Pk, folded k'-def]
 have k' \neq 0 x^{k'} = 1 by auto
 with not-less-Least[of - ?P, folded k'-def]
 have primitive-root-unity k' x **unfolding** primitive-root-unity-def by auto
 with primitive-root-unityD(3)[OF this assms]
 show ?thesis by auto
 qed

lemma primitive-root-unity-dvd: **fixes** x :: complex
assumes k: primitive-root-unity k x
 shows xⁿ = 1 \longleftrightarrow k dvd n
proof
 assume k dvd n **then obtain** j **where** n: n = k * j **unfolding** dvd-def by auto
 have xⁿ = (x^k)^j **unfolding** n power-mult by simp
 also have ... = 1 **unfolding** primitive-root-unityD[OF k] by simp
 finally show xⁿ = 1 .
next
 assume n: xⁿ = 1
 note k = primitive-root-unityD[OF k]
 show k dvd n
proof (cases n = 0)
 case n0: False
 from k(3)[OF n0] n **have** nk: n \geq k by force
 from roots-of-unity[OF k(1)] k(2) **obtain** i :: nat **where** xk: x = cis (i * 2 * pi / k)

```

    and ik: i < k by force
  from roots-of-unity[OF n0] n obtain j :: nat where xn: x = cis (j * 2 * pi /
n)
    and jn: j < n by force
  have cop: coprime i k
  proof (rule gcd-eq-1-imp-coprime)
    from k(1) have gcd i k ≠ 0 by auto
    from gcd-coprime-exists[OF this] this obtain i' k' g where
      *: i = i' * g k = k' * g g ≠ 0 and g: g = gcd i k by blast
    from *(2) k(1) have k': k' ≠ 0 by auto
    have x = cis (i * 2 * pi / k) by fact
    also have i * 2 * pi / k = i' * 2 * pi / k' unfolding * using *(3) by auto
    finally have x ^ k' = 1 by (simp add: DeMoivre k')
    with k(3)[OF k'] have k' ≥ k by linarith
    moreover with * k(1) have g = 1 by auto
    then show gcd i k = 1 by (simp add: g)
  qed
  from inj-onD[OF cis-inj-on xk[unfolded xn]] n0 k(1) ik jn
  have j * real k = i * real n by (auto simp: field-simps)
  hence real (j * k) = real (i * n) by simp
  hence eq: j * k = i * n by linarith
  with cop show k dvd n
    by (metis coprime-commute coprime-dvd-mult-right-iff dvd-triv-right)
  qed auto
  qed

```

lemma primitive-root-unity-simple-computation:

primitive-root-unity k x = (if k = 0 then False else

$x^k = 1 \wedge (\forall i \in \{1 ..< k\}. x^i \neq 1)$)

unfolding primitive-root-unity-def **by** auto

lemma primitive-root-unity-explicit: fixes x :: complex

shows primitive-root-unity 1 x \longleftrightarrow x = 1

primitive-root-unity 2 x \longleftrightarrow x = -1

primitive-root-unity 3 x \longleftrightarrow (x ∈ {Complex (-1/2) (sqrt 3 / 2), Complex (-1/2) (- sqrt 3 / 2)})

primitive-root-unity 4 x \longleftrightarrow (x ∈ {i, - i})

proof (atomize(full), goal-cases)

case 1

{

fix P :: nat \Rightarrow bool

have *: {1 ..< 2 :: nat} = {1} {1 ..< 3 :: nat} = {1,2} {1 ..< 4 :: nat} = {1,2,3}

by code-simp+

have ($\forall i \in \{1 ..< 2\}. P i$) = P 1 ($\forall i \in \{1 ..< 3\}. P i$) \longleftrightarrow P 1 \wedge P 2

($\forall i \in \{1 ..< 4\}. P i$) \longleftrightarrow P 1 \wedge P 2 \wedge P 3

unfolding * **by** auto

} note * = this

show ?case **unfolding** primitive-root-unity-simple-computation root-unity-explicit


```

*
  by (auto simp: complex-eq-iff)
qed

function decompose-prod-root-unity-main ::
  'a :: field poly  $\Rightarrow$  nat  $\Rightarrow$  nat list  $\times$  'a poly where
  decompose-prod-root-unity-main p k = (
    if k = 0 then ([], p) else
    let q = root-unity k in if q dvd p then if p = 0 then ([], 0) else
    map-prod (Cons k) id (decompose-prod-root-unity-main (p div q) k) else
    decompose-prod-root-unity-main p (k - 1))
  by pat-completeness auto

termination by (relation measure ( $\lambda$  (p,k). degree p + k), auto simp: degree-div-less)

declare decompose-prod-root-unity-main.simps[simp del]

lemma decompose-prod-root-unity-main: fixes p :: complex poly
  assumes p: p = prod-root-unity ks * f
  and d: decompose-prod-root-unity-main p k = (ks', g)
  and f:  $\bigwedge x. \text{cmod } x = 1 \implies \text{poly } f x \neq 0$ 
  and k:  $\bigwedge k'. k' > k \implies \neg \text{root-unity } k' \text{ dvd } p$ 
shows p = prod-root-unity ks' * f  $\wedge$  f = g  $\wedge$  set ks = set ks'
  using d p k
proof (induct p k arbitrary: ks ks' rule: decompose-prod-root-unity-main.induct)
  case (1 p k ks ks')
  note p = 1(4)
  note k = 1(5)
  from k[of Suc k] have p0: p  $\neq$  0 by auto
  hence p = 0  $\longleftrightarrow$  False by auto
  note d = 1(3)[unfolded decompose-prod-root-unity-main.simps[of p k] this if-False]
  Let-def]
  from p0[unfolded p] have ks0: 0  $\notin$  set ks by simp
  from f[of 1] have f0: f  $\neq$  0 by auto
  note IH = 1(1)[OF - refl - p0] 1(2)[OF - refl]
  show ?case
  proof (cases k = 0)
    case True
    with p k[unfolded this, of hd ks] p0 have ks = []
      by (cases ks, auto simp: prod-root-unity-def)
    with d p True show ?thesis by (auto simp: prod-root-unity-def)
  next
    case k0: False
    note IH = IH[OF k0]
    from k0 have k = 0  $\longleftrightarrow$  False by auto
    note d = d[unfolded this if-False]
    let ?u = root-unity k :: complex poly
    show ?thesis
    proof (cases ?u dvd p)

```

```

case True
note IH = IH(1)[OF True]
let ?call = decompose-prod-root-unity-main (p div ?u) k
from True d obtain Ks where rec: ?call = (Ks,g) and ks': ks' = (k # Ks)
  by (cases ?call, auto)
from True have ?u dvd p  $\longleftrightarrow$  True by simp
note d = d[unfolded this if-True rec]
let ?x = cis (2 * pi / k)
have rt: poly ?u ?x = 0 unfolding poly-root-unity using cis-times-2pi[of 1]
  by (simp add: DeMoiivre)
with True have poly p ?x = 0 unfolding dvd-def by auto
from this[unfolded p] f[of ?x] rt have poly (prod-root-unity ks) ?x = 0
  unfolding poly-root-unity by auto
from this[unfolded poly-prod-root-unity] ks0 obtain k' where k': k'  $\in$  set ks
  and rt: ?x ^ k' = 1 and k'0: k'  $\neq$  0 by auto
let ?u' = root-unity k' :: complex poly
from k' rt k'0 have rtk': poly ?u' ?x = 0 unfolding poly-root-unity by auto
{
  let ?phi = k' * (2 * pi / k)
  assume k' < k
  hence 0 < ?phi ?phi < 2 * pi using k0 k'0 by (auto simp: field-simps)
  from cis-plus-2pi-neq-1[OF this] rtk'
  have False unfolding poly-root-unity DeMoiivre ..
}
hence kk': k  $\leq$  k' by presburger
{
  assume k' > k
  from k[OF this, unfolded p]
  have  $\neg$  ?u' dvd prod-root-unity ks using dvd-mult2 by auto
  with k' have False unfolding prod-root-unity-def
    using prod-list-dvd[of ?u' map root-unity ks] by auto
}
with kk' have kk': k' = k by presburger
with k' have k  $\in$  set ks by auto
from split-list[OF this] obtain ks1 ks2 where ks: ks = ks1 @ k # ks2 by
auto
hence p div ?u = (?u * (prod-root-unity (ks1 @ ks2) * f)) div ?u
  by (simp add: ac-simps p prod-root-unity-def)
also have ... = prod-root-unity (ks1 @ ks2) * f
  by (rule nonzero-mult-div-cancel-left, insert k0, auto)
finally have id: p div ?u = prod-root-unity (ks1 @ ks2) * f .
from d have ks': ks' = k # Ks by auto
have k < k'  $\implies$   $\neg$  root-unity k' dvd p div ?u for k'
  using k[of k'] True by (metis dvd-div-mult-self dvd-mult2)
from IH[OF rec id this]
have id: p div root-unity k = prod-root-unity Ks * f and
  *: f = g  $\wedge$  set (ks1 @ ks2) = set Ks by auto
from arg-cong[OF id, of  $\lambda x. x * ?u$ ] True
have p = prod-root-unity Ks * f * root-unity k by auto

```

```

    thus ?thesis using * unfolding ks ks' by (auto simp: prod-root-unity-def)
  next
    case False
    from d False have decompose-prod-root-unity-main p (k - 1) = (ks',g) by
  auto
    note IH = IH(2)[OF False this p]
    have k: k - 1 < k'  $\implies$   $\neg$  root-unity k' dvd p for k' using False k[of k'] k0
      by (cases k' = k, auto)
    show ?thesis by (rule IH, insert False k, auto)
  qed
qed
qed

```

definition *decompose-prod-root-unity* p = *decompose-prod-root-unity-main* p (degree p)

lemma *decompose-prod-root-unity*: fixes p :: complex poly
 assumes p: p = prod-root-unity ks * f
 and d: decompose-prod-root-unity p = (ks',g)
 and f: $\bigwedge x. \text{cmod } x = 1 \implies \text{poly } f x \neq 0$
 and p0: p \neq 0
 shows p = prod-root-unity ks' * f \wedge f = g \wedge set ks = set ks'
proof (rule decompose-prod-root-unity-main[OF p d[unfolded decompose-prod-root-unity-def] f])
 fix k
 assume deg: degree p < k
 hence degree p < degree (root-unity k) by simp
 with p0 show \neg root-unity k dvd p
 by (simp add: poly-divides-conv0)
 qed

lemma (in comm-ring-hom) hom-root-unity: map-poly hom (root-unity n) = root-unity n
proof –
 interpret p: map-poly-comm-ring-hom hom ..
 show ?thesis unfolding root-unity-def
 by (simp add: hom-distrib)
 qed

lemma (in idom-hom) hom-prod-root-unity: map-poly hom (prod-root-unity n) = prod-root-unity n
proof –
 interpret p: map-poly-comm-ring-hom hom ..
 show ?thesis unfolding prod-root-unity-def p.hom-prod-list map-map o-def hom-root-unity ..
 qed

lemma (in field-hom) hom-decompose-prod-root-unity-main:
 decompose-prod-root-unity-main (map-poly hom p) k = map-prod id (map-poly

```

hom)
  (decompose-prod-root-unity-main p k)
proof (induct p k rule: decompose-prod-root-unity-main.induct)
  case (1 p k)
  let ?h = map-poly hom
  let ?p = ?h p
  let ?u = root-unity k :: 'a poly
  let ?u' = root-unity k :: 'b poly
  interpret p: map-poly-inj-idom-divide-hom hom ..
  have u': ?u' = ?h ?u unfolding hom-root-unity ..
  note simp = decompose-prod-root-unity-main.simps
  let ?rec1 = decompose-prod-root-unity-main (p div ?u) k
  have 0: ?p = 0  $\longleftrightarrow$  p = 0 by simp
  show ?case
    unfolding simp[of ?p k] simp[of p k] if-distrib[of map-prod id ?h] Let-def u'
    unfolding 0 p.hom-div[symmetric] p.hom-dvd-iff
    by (rule if-cong[OF refl], force, rule if-cong[OF refl if-cong[OF refl]], force,
      (subst 1(1), auto, cases ?rec1, auto)[1],
      (subst 1(2), auto))
qed

```

```

lemma (in field-hom) hom-decompose-prod-root-unity:
  decompose-prod-root-unity (map-poly hom p) = map-prod id (map-poly hom)
  (decompose-prod-root-unity p)
  unfolding decompose-prod-root-unity-def
  by (subst hom-decompose-prod-root-unity-main, simp)

```

end

5.1 The Perron Frobenius Theorem for Irreducible Matrices

```

theory Perron-Frobenius-Irreducible

```

```

imports

```

```

  Perron-Frobenius

```

```

  Roots-Unity

```

```

  Rank-Nullity-Theorem.Miscellaneous

```

```

begin

```

```

lifting-forget vec.lifting

```

```

lifting-forget mat.lifting

```

```

lifting-forget poly.lifting

```

```

lemma charpoly-of-real: charpoly (map-matrix complex-of-real A) = map-poly of-real
(charpoly A)

```

```

  by (transfer-hma rule: of-real-hom.char-poly-hom)

```

```

context includes lifting-syntax

```

```

begin

```

lemma *HMA-M-smult*[*transfer-rule*]: $((=) \implies \text{HMA-M} \implies \text{HMA-M}) (\cdot_m)$
 $((*k))$
unfolding *smult-mat-def*
unfolding *rel-fun-def HMA-M-def from-hma_m-def*
by (*auto simp: matrix-scalar-mult-def*)
end

lemma *order-charpoly-smult*: **fixes** $A :: \text{complex } ^\wedge 'n \ ^\wedge 'n$
assumes $k: k \neq 0$
shows $\text{order } x (\text{charpoly } (k *k A)) = \text{order } (x / k) (\text{charpoly } A)$
by (*transfer fixing: k, rule order-char-poly-smult[OF - k]*)

lemma *smult-eigen-vector*: **fixes** $a :: 'a :: \text{field}$
assumes *eigen-vector A v x*
shows *eigen-vector (a *k A) v (a * x)*
proof –
from *assms[unfolded eigen-vector-def]* **have** $v: v \neq 0$ **and** *id: A *v v = x *s v*
by *auto*
from *arg-cong[OF id, of (*s) a]* **have** *id: (a *k A) *v v = (a * x) *s v*
unfolding *scalar-matrix-vector-assoc* **by** *simp*
thus *eigen-vector (a *k A) v (a * x)* **using** v **unfolding** *eigen-vector-def* **by**
auto
qed

lemma *smult-eigen-value*: **fixes** $a :: 'a :: \text{field}$
assumes *eigen-value A x*
shows *eigen-value (a *k A) (a * x)*
using *assms smult-eigen-vector[of A - x a]* **unfolding** *eigen-value-def* **by** *blast*

locale *fixed-mat* = **fixes** $A :: 'a :: \text{zero } ^\wedge 'n \ ^\wedge 'n$
begin
definition $G :: 'n \text{ rel}$ **where**
 $G = \{ (i,j). A \$ i \$ j \neq 0 \}$

definition *irreducible* :: *bool* **where**
 $\text{irreducible} = (\text{UNIV} \subseteq G^{\wedge+})$
end

lemma *G-transpose*:
 $\text{fixed-mat}.G (\text{transpose } A) = ((\text{fixed-mat}.G A))^{\wedge-1}$
unfolding *fixed-mat.G-def* **by** (*force simp: transpose-def*)

lemma *G-transpose-trancl*:
 $(\text{fixed-mat}.G (\text{transpose } A))^{\wedge+} = ((\text{fixed-mat}.G A)^{\wedge+})^{\wedge-1}$
unfolding *G-transpose trancl-converse* **by** *auto*

locale *pf-nonneg-mat* = *fixed-mat A* **for**
 $A :: 'a :: \text{linordered-idom } ^\wedge 'n \ ^\wedge 'n +$

```

    assumes non-neg-mat: non-neg-mat A
  begin
  lemma nonneg:  $A \ $ i \ $ j \ge 0$ 
    using non-neg-mat unfolding non-neg-mat-def elements-mat-h-def by auto

  lemma nonneg-matpow:  $\text{matpow } A \ n \ \$ i \ \$ j \ge 0$ 
    by (induct n arbitrary: i j, insert nonneg,
        auto intro!: sum-nonneg simp: matrix-matrix-mult-def mat-def)

  lemma G-relpow-matpow-pos:  $(i,j) \in G \ \sim n \implies \text{matpow } A \ n \ \$ i \ \$ j > 0$ 
  proof (induct n arbitrary: i j)
    case (0 i)
    thus ?case by (auto simp: mat-def)
  next
    case (Suc n i j)
    from Suc(2) have  $(i,j) \in G \ \sim n \ O \ G$ 
      by (simp add: relpow-commute)
    then obtain k where
      ik:  $A \ \$ k \ \$ j \neq 0$  and  $kj: (i, k) \in G \ \sim n$  by (auto simp: G-def)
    from ik nonneg[of k j] have  $ik: A \ \$ k \ \$ j > 0$  by auto
    from Suc(1)[OF kj] have IH:  $\text{matpow } A \ n \ \$ i \ \$ k > 0$  .
    thus ?case using ik by (auto simp: nonneg-matpow nonneg matrix-matrix-mult-def

      intro!: sum-pos2[of - k] mult-nonneg-nonneg)
  qed

  lemma matpow-mono: assumes  $B: \bigwedge i j. B \ \$ i \ \$ j \ge A \ \$ i \ \$ j$ 
    shows  $\text{matpow } B \ n \ \$ i \ \$ j \ge \text{matpow } A \ n \ \$ i \ \$ j$ 
  proof (induct n arbitrary: i j)
    case (Suc n i j)
    thus ?case using B nonneg-matpow[of n] nonneg
      by (auto simp: matrix-matrix-mult-def intro!: sum-mono mult-mono')
  qed simp

  lemma matpow-sum-one-mono:  $\text{matpow } (A + \text{mat } 1) \ (n + k) \ \$ i \ \$ j \ge \text{matpow}$ 
 $(A + \text{mat } 1) \ n \ \$ i \ \$ j$ 
  proof (induct k)
    case (Suc k)
    have  $(\text{matpow } (A + \text{mat } 1) \ (n + k) \ ** \ A) \ \$ i \ \$ j \ge 0$  unfolding ma-
      trix-matrix-mult-def
      using order.trans[OF nonneg-matpow matpow-mono[of A + mat 1 n + k]]
      by (auto intro!: sum-nonneg mult-nonneg-nonneg nonneg simp: mat-def)
    thus ?case using Suc by (simp add: matrix-add-ldistrib matrix-mul-rid)
  qed simp

  lemma G-relpow-matpow-pos-ge:
    assumes  $(i,j) \in G \ \sim m \ n \ge m$ 
    shows  $\text{matpow } (A + \text{mat } 1) \ n \ \$ i \ \$ j > 0$ 
  proof -

```

from *assms(2)* **obtain** *k* **where** $n: n = m + k$ **using** *le-Suc-ex* **by** *blast*
have $0 < \text{matpow } A \ m \ \$ \ i \ \$ \ j$ **by** (rule *G-relpow-matpow-pos[OF assms(1)]*)
also have $\dots \leq \text{matpow } (A + \text{mat } 1) \ m \ \$ \ i \ \$ \ j$
by (rule *matpow-mono, auto simp: mat-def*)
also have $\dots \leq \text{matpow } (A + \text{mat } 1) \ n \ \$ \ i \ \$ \ j$ **unfolding** *n* **using** *mat-pow-sum-one-mono* .
finally show *?thesis* .
qed
end

locale *perron-frobenius* = *pf-nonneg-mat A*
for $A :: \text{real}^{\wedge 'n} \wedge 'n +$
assumes *irr: irreducible*
begin

definition *N* **where** $N = (\text{SOME } N. \forall ij. \exists n \leq N. ij \in G^{\wedge n})$

lemma *N: $\exists n \leq N. ij \in G^{\wedge n}$*
proof –
{
fix *ij*
have $ij \in G^+$ **using** *irr[unfolded irreducible-def]* **by** *auto*
from *this[unfolded tranc1-power]* **have** $\exists n. ij \in G^{\wedge n}$ **by** *blast*
}
hence $\forall ij. \exists n. ij \in G^{\wedge n}$ **by** *auto*
from *choice[OF this]* **obtain** *f* **where** $f: \bigwedge ij. ij \in G^{\wedge (f \ ij)}$ **by** *auto*
define *N* **where** $N: N = \text{Max } (\text{range } f)$
{
fix *ij*
from *f[of ij]* **have** $ij \in G^{\wedge (f \ ij)}$.
moreover have $f \ ij \leq N$ **unfolding** *N*
by (rule *Max-ge, auto*)
ultimately have $\exists n \leq N. ij \in G^{\wedge n}$ **by** *blast*
}
note *main = this*
let $?P = \lambda N. \forall ij. \exists n \leq N. ij \in G^{\wedge n}$
from *main* **have** $?P \ N$ **by** *blast*
from *someI[of ?P, OF this, folded N-def]*
show *?thesis* **by** *blast*
qed

lemma *irreducible-matpow-pos: assumes irreducible*
shows $\text{matpow } (A + \text{mat } 1) \ N \ \$ \ i \ \$ \ j > 0$
proof –
from *N* **obtain** *n* **where** $n: n \leq N$ **and** *reach: $(i,j) \in G^{\wedge n}$* **by** *auto*
show *?thesis* **by** (rule *G-relpow-matpow-pos-ge[OF reach n]*)
qed

lemma *pf-transpose: perron-frobenius (transpose A)*
proof

show *fixed-mat.irreducible* (*transpose A*)
unfolding *fixed-mat.irreducible-def G-transpose-trancl* **using** *irr[unfolded irreducible-def]*
by *auto*
qed (*insert nonneg, auto simp: transpose-def non-neg-mat-def elements-mat-h-def*)

abbreviation *le-vec* :: $\text{real}^n \Rightarrow \text{real}^n \Rightarrow \text{bool}$ **where**
le-vec $x\ y \equiv (\forall\ i.\ x\ \$\ i \leq y\ \$\ i)$

abbreviation *lt-vec* :: $\text{real}^n \Rightarrow \text{real}^n \Rightarrow \text{bool}$ **where**
lt-vec $x\ y \equiv (\forall\ i.\ x\ \$\ i < y\ \$\ i)$

definition $A1n = \text{matpow}\ (A + \text{mat}\ 1)\ N$

lemmas $A1n\text{-pos} = \text{irreducible-matpow-pos}[OF\ \text{irr},\ \text{folded}\ A1n\text{-def}]$

definition *r* :: $\text{real}^n \Rightarrow \text{real}$ **where**
r $x = \text{Min}\ \{ (A * v\ x)\ \$\ j / x\ \$\ j \mid j.\ x\ \$\ j \neq 0 \}$

definition *X* :: $(\text{real}^n)\text{set}$ **where**
X = $\{ x.\ \text{le-vec}\ 0\ x \wedge x \neq 0 \}$

lemma *nonneg-Ax*: $x \in X \implies \text{le-vec}\ 0\ (A * v\ x)$
unfolding *X-def* **using** *nonneg*
by (*auto simp: matrix-vector-mult-def intro!: sum-nonneg*)

lemma *A-nonzero-fixed-i*: $\exists\ j.\ A\ \$\ i\ \$\ j \neq 0$
proof –
from *irr[unfolded irreducible-def]* **have** $(i,i) \in G^+$ **by** *auto*
then obtain *j* **where** $(i,j) \in G$ **by** (*metis converse-tranclE*)
hence *Aij*: $A\ \$\ i\ \$\ j \neq 0$ **unfolding** *G-def* **by** *auto*
thus *?thesis ..*
qed

lemma *A-nonzero-fixed-j*: $\exists\ i.\ A\ \$\ i\ \$\ j \neq 0$
proof –
from *irr[unfolded irreducible-def]* **have** $(j,j) \in G^+$ **by** *auto*
then obtain *i* **where** $(i,j) \in G$ **by** (*cases, auto*)
hence *Aij*: $A\ \$\ i\ \$\ j \neq 0$ **unfolding** *G-def* **by** *auto*
thus *?thesis ..*
qed

lemma *Ax-pos*: **assumes** $x:\ \text{lt-vec}\ 0\ x$
shows $\text{lt-vec}\ 0\ (A * v\ x)$
proof
fix *i*
from *A-nonzero-fixed-i[of i]* **obtain** *j* **where** $A\ \$\ i\ \$\ j \neq 0$ **by** *auto*
with *nonneg[of i j]* **have** $A:\ A\ \$\ i\ \$\ j > 0$ **by** *simp*
from *x* **have** $x\ \$\ j \geq 0$ **for** *j* **by** (*auto simp: order.strict-iff-order*)

note $\text{nonneg} = \text{mult-nonneg-nonneg}[OF \text{nonneg}[of i] \text{this}]$
have $(A *v x) \$ i = (\sum_{j \in UNIV}. A \$ i \$ j * x \$ j)$
unfolding $\text{matrix-vector-mult-def}$ **by** simp
also have $\dots = A \$ i \$ j * x \$ j + (\sum_{j \in UNIV - \{j\}}. A \$ i \$ j * x \$ j)$
by $(\text{subst sum.remove}, \text{auto})$
also have $\dots > 0 + 0$
by $(\text{rule add-less-le-mono}, \text{insert } A \ x[\text{rule-format}] \ \text{nonneg},$
 $\text{auto intro!}: \text{sum-nonneg mult-pos-pos})$
finally show $0 \$ i < (A *v x) \$ i$ **by** simp
qed

lemma nonzero-Ax: assumes $x: x \in X$
shows $A *v x \neq 0$

proof

assume $0: A *v x = 0$
from $x[\text{unfolded } X\text{-def}]$ **have** $x: \text{le-vec } 0 \ x \ x \neq 0$ **by** auto
from $x(2)$ **obtain** j **where** $xj: x \$ j \neq 0$
by $(\text{metis vec-eq-iff zero-index})$
from $A\text{-nonzero-fixed-}j[\text{of } j]$ **obtain** i **where** $Aij: A \$ i \$ j \neq 0$ **by** auto
from $\text{arg-cong}[OF 0, \text{of } \lambda v. v \$ i, \text{unfolded matrix-vector-mult-def}]$
have $0 = (\sum_{k \in UNIV}. A \$ h \ i \ \$ h \ k * x \$ h \ k)$ **by** auto
also have $\dots = A \$ h \ i \ \$ h \ j * x \$ h \ j + (\sum_{k \in UNIV - \{j\}}. A \$ h \ i \ \$ h \ k * x$
 $\$ h \ k)$
by $(\text{subst sum.remove}[of - j], \text{auto})$
also have $\dots > 0 + 0$
by $(\text{rule add-less-le-mono}, \text{insert } \text{nonneg}[of i] \ Aij \ x(1) \ xj,$
 $\text{auto intro!}: \text{sum-nonneg mult-pos-pos simp: dual-order.not-eq-order-implies-strict})$

finally show False **by** simp

qed

lemma r-witness: assumes $x: x \in X$

shows $\exists j. x \$ j > 0 \wedge r x = (A *v x) \$ j / x \$ j$

proof –

from $x[\text{unfolded } X\text{-def}]$ **have** $x: \text{le-vec } 0 \ x \ x \neq 0$ **by** auto
let $?A = \{ (A *v x) \$ j / x \$ j \mid j. x \$ j \neq 0 \}$
from $x(2)$ **obtain** j **where** $x \$ j \neq 0$
by $(\text{metis vec-eq-iff zero-index})$
hence $\text{empty}: ?A \neq \{\}$ **by** auto
from $\text{Min-in}[OF - \text{this}, \text{folded } r\text{-def}]$
obtain j **where** $x \$ j \neq 0$ **and** $rx: r x = (A *v x) \$ j / x \$ j$ **by** auto
with x **have** $x \$ j > 0$ **by** $(\text{auto simp: dual-order.not-eq-order-implies-strict})$
with rx **show** $?thesis$ **by** auto

qed

lemma rx-nonneg: assumes $x: x \in X$

shows $r x \geq 0$

proof –

```

from  $x$ [unfolded X-def] have  $x$ : le-vec 0  $x$   $x \neq 0$  by auto
let  $?A = \{ (A *v x) \$ j / x \$ j \mid j. x \$ j \neq 0 \}$ 
from r-witness[OF  $\langle x \in X \rangle$ ]
have empty:  $?A \neq \{\}$  by force
show ?thesis unfolding r-def X-def
proof (subst Min-ge-iff, force, use empty in force, intro ballI)
  fix  $y$ 
  assume  $y \in ?A$ 
  then obtain  $j$  where  $y = (A *v x) \$ j / x \$ j$  and  $x \$ j \neq 0$  by auto
  from nonneg-Ax[OF  $\langle x \in X \rangle$ ] this  $x$ 
  show  $0 \leq y$  by simp
qed
qed

lemma rx-pos: assumes  $x$ : lt-vec 0  $x$ 
shows  $r x > 0$ 
proof –
  from Ax-pos[OF  $x$ ] have lt: lt-vec 0  $(A *v x)$  .
  from  $x$  have  $x'$ :  $x \in X$  unfolding X-def order.strict-iff-order by auto
  let  $?A = \{ (A *v x) \$ j / x \$ j \mid j. x \$ j \neq 0 \}$ 
  from r-witness[OF  $\langle x \in X \rangle$ ]
  have empty:  $?A \neq \{\}$  by force
  show ?thesis unfolding r-def X-def
  proof (subst Min-gr-iff, force, use empty in force, intro ballI)
    fix  $y$ 
    assume  $y \in ?A$ 
    then obtain  $j$  where  $y = (A *v x) \$ j / x \$ j$  and  $x \$ j \neq 0$  by auto
    from lt this  $x$  show  $0 < y$  by simp
  qed
qed

lemma rx-le-Ax: assumes  $x$ :  $x \in X$ 
shows le-vec  $(r x *s x) (A *v x)$ 
proof (intro allI)
  fix  $i$ 
  show  $(r x *s x) \$h i \leq (A *v x) \$h i$ 
  proof (cases  $x \$ i = 0$ )
    case True
    with nonneg-Ax[OF  $x$ ] show ?thesis by auto
  next
  case False
  with  $x$ [unfolded X-def] have pos:  $x \$ i > 0$ 
  by (auto simp: dual-order.not-eq-order-implies-strict)
  from False have  $(A *v x) \$h i / x \$ i \in \{ (A *v x) \$ j / x \$ j \mid j. x \$ j \neq 0 \}$ 
by auto
  hence  $(A *v x) \$h i / x \$ i \geq r x$  unfolding r-def by simp
  hence  $x \$ i * r x \leq x \$ i * ((A *v x) \$h i / x \$ i)$  unfolding mult-le-cancel-left-pos[OF pos] .
  also have  $\dots = (A *v x) \$h i$  using pos by simp

```

finally show *?thesis* **by** (*simp add: ac-simps*)
qed
qed

lemma *rho-le-x-Ax-imp-rho-le-rx*: **assumes** $x: x \in X$
and $\varrho: \text{le-vec } (\varrho *s x) (A *v x)$
shows $\varrho \leq r x$
proof –
from *r-witness*[*OF x*] **obtain** j **where**
 $rx: r x = (A *v x) \$ j / x \$ j$ **and** $xj: x \$ j > 0 \ x \$ j \geq 0$ **by** *auto*
from *divide-right-mono*[*OF* ϱ [*rule-format, of j*] $xj(2)$]
show *?thesis* **unfolding** rx **using** xj **by** *simp*
qed

lemma *rx-Max*: **assumes** $x: x \in X$
shows $r x = \text{Sup } \{ \varrho . \text{le-vec } (\varrho *s x) (A *v x) \}$ (**is** $- = \text{Sup } ?S$)
proof –
have $r x \in ?S$ **using** *rx-le-Ax*[*OF x*] **by** *auto*
moreover {
fix y
assume $y \in ?S$
hence $y: \text{le-vec } (y *s x) (A *v x)$ **by** *auto*
from *rho-le-x-Ax-imp-rho-le-rx*[*OF x this*]
have $y \leq r x$.
}
ultimately show *?thesis* **by** (*metis (mono-tags, lifting) cSup-eq-maximum*)
qed

lemma *r-smult*: **assumes** $x: x \in X$
and $a: a > 0$
shows $r (a *s x) = r x$
unfolding *r-def*
by (*rule arg-cong*[*of - - Min*], *unfold vector-smult-distrib, insert a, simp*)

definition $X1 = (X \cap \{x. \text{norm } x = 1\})$

lemma *bounded-X1*: *bounded X1* **unfolding** *bounded-iff X1-def* **by** *auto*

lemma *closed-X1*: *closed X1*

proof –
have $X1: X1 = \{x. \text{le-vec } 0 x \wedge \text{norm } x = 1\}$
unfolding *X1-def X-def* **by** *auto*
show *?thesis* **unfolding** $X1$
by (*intro closed-Collect-conj closed-Collect-all closed-Collect-le closed-Collect-eq,*
auto intro: continuous-intros)
qed

lemma *compact-X1*: *compact X1* **using** *bounded-X1 closed-X1*
by (*simp add: compact-eq-bounded-closed*)

definition $\text{pow-A-1 } x = A1n *v x$

lemma $\text{continuous-pow-A-1: continuous-on } R \text{ pow-A-1}$
unfolding $\text{pow-A-1-def continuous-on}$
by $(\text{auto intro: tendsto-intros})$

definition $Y = \text{pow-A-1 } ' X1$

lemma $\text{compact-Y: compact } Y$
unfolding $Y\text{-def using compact-X1 continuous-pow-A-1[of X1]$
by $(\text{metis compact-continuous-image})$

lemma $Y\text{-pos-main: assumes } y: y \in \text{pow-A-1 } ' X$
shows $y \ \$ \ i > 0$

proof $-$

from y **obtain** x **where** $x: x \in X$ **and** $y: y = \text{pow-A-1 } x$ **unfolding** $Y\text{-def}$
 $X1\text{-def}$ **by** auto

from $r\text{-witness}[OF x]$ **obtain** j **where** $xj: x \ \$ \ j > 0$ **by** auto

from $x[\text{unfolded } X\text{-def}]$ **have** $xi: x \ \$ \ i \geq 0$ **for** i **by** auto

have $\text{nonneg: } 0 \leq A1n \ \$ \ i \ \$ \ k * x \ \$ \ k$ **for** k **using** $A1n\text{-pos}[of i k] xi[of k]$ **by**
 auto

have $y \ \$ \ i = (\sum j \in UNIV. A1n \ \$ \ i \ \$ \ j * x \ \$ \ j)$

unfolding $y \ \text{pow-A-1-def matrix-vector-mult-def}$ **by** simp

also have $\dots = A1n \ \$ \ i \ \$ \ j * x \ \$ \ j + (\sum j \in UNIV - \{j\}. A1n \ \$ \ i \ \$ \ j * x \ \$ \ j)$

by $(\text{subst sum.remove, auto})$

also have $\dots > 0 + 0$

by $(\text{rule add-less-le-mono, insert } xj \ A1n\text{-pos nonneg,}$

$\text{auto intro!: sum-nonneg mult-pos-pos simp: dual-order.not-eq-order-implies-strict})$

finally show $?thesis$ **by** simp

qed

lemma $Y\text{-pos: assumes } y: y \in Y$

shows $y \ \$ \ i > 0$

using $Y\text{-pos-main}[of y i] y$ **unfolding** $Y\text{-def } X1\text{-def}$ **by** auto

lemma $Y\text{-nonzero: assumes } y: y \in Y$

shows $y \ \$ \ i \neq 0$

using $Y\text{-pos}[OF y, of i]$ **by** auto

definition $r' :: \text{real} \wedge 'n \Rightarrow \text{real}$ **where**

$r' x = \text{Min } (\text{range } (\lambda j. (A *v x) \ \$ \ j / x \ \$ \ j))$

lemma $r'\text{-r: assumes } x: x \in Y$ **shows** $r' x = r x$

unfolding $r'\text{-def } r\text{-def}$

proof $(\text{rule arg-cong}[of - - Min])$

have $\text{range } (\lambda j. (A *v x) \ \$ \ j / x \ \$ \ j) \subseteq \{(A *v x) \ \$ \ j / x \ \$ \ j \mid j. x \ \$ \ j \neq 0\}$ **(is**

$?L \subseteq ?R$)
proof
 fix y
 assume $y \in ?L$
 then obtain j **where** $y = (A *v x) \$ j / x \$ j$ **by** *auto*
 with $Y\text{-pos}[OF\ x, of\ j]$ **show** $y \in ?R$ **by** *auto*
qed
moreover have $?L \supseteq ?R$ **by** *auto*
ultimately show $?L = ?R$ **by** *blast*
qed

lemma *continuous-Y-r: continuous-on Y r*
proof –
 have $*$: $(\forall y \in Y. P\ y\ (r\ y)) = (\forall y \in Y. P\ y\ (r'\ y))$ **for** P **using** $r'\text{-}r$ **by** *auto*
 have $\text{continuous-on } Y\ r = \text{continuous-on } Y\ r'$
 by $(\text{rule } \text{continuous-on-cong}[OF\ refl\ r'\text{-}r[\text{symmetric}]])$
 also have ...
 unfolding $\text{continuous-on } r'\text{-}def$ **using** $Y\text{-nonzero}$
 by $(\text{auto } \text{intro!}: \text{tendsto-Min } \text{tendsto-intros})$
 finally show $?thesis$.
qed

lemma $X1\text{-nonempty}: X1 \neq \{\}$
proof –
 define x **where** $x = ((\chi\ i. \text{if } i = \text{undefined then } 1 \text{ else } 0) :: \text{real } ^n)$
 {
 assume $x = 0$
 from $\text{arg-cong}[OF\ \text{this}, of\ \lambda\ x. x\ \$\ \text{undefined}]$ **have** False **unfolding** $x\text{-def}$ **by** *auto*
 }
 hence $x: x \neq 0$ **by** *auto*
 moreover have $\text{le-vec } 0\ x$ **unfolding** $x\text{-def}$ **by** *auto*
 moreover have $\text{norm } x = 1$ **unfolding** $\text{norm-vec-def } L2\text{-set-def}$
 by $(\text{auto}, \text{subst } \text{sum.remove}[of\ \text{-undefined}], \text{auto } \text{simp}: x\text{-def})$
 ultimately show $?thesis$ **unfolding** $X1\text{-def } X\text{-def}$ **by** *auto*
qed

lemma $Y\text{-nonempty}: Y \neq \{\}$
 unfolding $Y\text{-def}$ **using** $X1\text{-nonempty}$ **by** *auto*

definition z **where** $z = (\text{SOME } z. z \in Y \wedge (\forall y \in Y. r\ y \leq r\ z))$

abbreviation $sr \equiv r\ z$

lemma $z: z \in Y$ **and** $sr\text{-max-}Y: \bigwedge y. y \in Y \implies r\ y \leq sr$

proof –
 let $?P = \lambda\ z. z \in Y \wedge (\forall y \in Y. r\ y \leq r\ z)$
 from $\text{continuous-attains-sup}[OF\ \text{compact-}Y\ Y\text{-nonempty } \text{continuous-}Y\text{-}r]$
 obtain y **where** $?P\ y$ **by** *blast*

from *someI*[*of ?P, OF this, folded z-def*]
show $z \in Y \wedge y. y \in Y \implies r y \leq r z$ **by** *blast+*
qed

lemma *Y-subset-X*: $Y \subseteq X$
proof
fix *y*
assume $y \in Y$
from *Y-pos*[*OF this*] **show** $y \in X$ **unfolding** *X-def*
by (*auto simp: order.strict-iff-order*)
qed

lemma *zX*: $z \in X$
using *z(1) Y-subset-X* **by** *auto*

lemma *le-vec-mono-left*: **assumes** $B: \bigwedge i j. B \$ i \$ j \geq 0$
and *le-vec x y*
shows *le-vec* ($B *v x$) ($B *v y$)
proof (*intro allI*)
fix *i*
show ($B *v x$) $\$ i \leq (B *v y) \$ i$ **unfolding** *matrix-vector-mult-def* **using** *B*[*of i*]
assms(2)
by (*auto intro!: sum-mono mult-left-mono*)
qed

lemma *matpow-1-commute*: $\text{matpow } (A + \text{mat } 1) n ** A = A ** \text{matpow } (A + \text{mat } 1) n$
by (*induct n, auto simp: matrix-add-rdistrib matrix-add-ldistrib matrix-mul-rid matrix-mul-lid matrix-mul-assoc[symmetric]*)

lemma *A1n-commute*: $A1n ** A = A ** A1n$
unfolding *A1n-def* **by** (*rule matpow-1-commute*)

lemma *le-vec-pow-A-1*: **assumes** *le*: *le-vec* ($\rho *s x$) ($A *v x$)
shows *le-vec* ($\rho *s \text{pow-A-1 } x$) ($A *v \text{pow-A-1 } x$)
proof –
have $A1n \$ i \$ j \geq 0$ **for** $i j$ **using** *A1n-pos*[*of i j*] **by** *auto*
from *le-vec-mono-left*[*OF this le*]
have *le-vec* ($A1n *v (\rho *s x)$) ($A1n *v (A *v x)$) .
also have $A1n *v (A *v x) = (A1n ** A) *v x$ **by** (*simp add: matrix-vector-mul-assoc*)
also have $\dots = A *v (A1n *v x)$ **unfolding** *A1n-commute* **by** (*simp add: matrix-vector-mul-assoc*)
also have $\dots = A *v (\text{pow-A-1 } x)$ **unfolding** *pow-A-1-def* ..
also have $A1n *v (\rho *s x) = \rho *s (A1n *v x)$ **unfolding** *vector-smult-distrib*
..
also have $\dots = \rho *s \text{pow-A-1 } x$ **unfolding** *pow-A-1-def* ..
finally show *le-vec* ($\rho *s \text{pow-A-1 } x$) ($A *v \text{pow-A-1 } x$) .

qed

lemma *r-pow-A-1*: **assumes** $x: x \in X$

shows $r x \leq r (pow-A-1 x)$

proof –

let $?y = pow-A-1 x$

have $?y \in pow-A-1 ' X$ **using** x **by** *auto*

from *Y-pos-main*[*OF this*]

have $y: ?y \in X$ **unfolding** *X-def* **by** (*auto simp: order.strict-iff-order*)

let $?A = \{\varrho. le-vec (\varrho *s x) (A *v x)\}$

let $?B = \{\varrho. le-vec (\varrho *s pow-A-1 x) (A *v pow-A-1 x)\}$

show *?thesis* **unfolding** *rx-Max*[*OF x*] *rx-Max*[*OF y*]

proof (*rule cSup-mono*)

show *bdd-above* $?B$ **using** *rho-le-x-Ax-imp-rho-le-rx*[*OF y*] **by** *fast*

show $?A \neq \{\}$ **using** *rx-le-Ax*[*OF x*] **by** *auto*

fix ρ

assume $\rho \in ?A$

hence *le-vec* ($\rho *s x$) ($A *v x$) **by** *auto*

from *le-vec-pow-A-1*[*OF this*] **have** $\rho \in ?B$ **by** *auto*

thus $\exists \rho' \in ?B. \rho \leq \rho'$ **by** *auto*

qed

qed

lemma *sr-max*: **assumes** $x: x \in X$

shows $r x \leq sr$

proof –

let $?n = norm x$

define x' **where** $x' = inverse ?n *s x$

from x [*unfolded X-def*] **have** $x0: x \neq 0$ **by** *auto*

hence $n: ?n > 0$ **by** *auto*

have $x': x' \in X1$ $x' \in X$ **using** x n **unfolding** *X1-def* *X-def* *x'-def* **by** (*auto simp: norm-smult*)

have *id*: $r x = r x'$ **unfolding** *x'-def*

by (*rule sym, rule r-smult*[*OF x*], *insert n, auto*)

define y **where** $y = pow-A-1 x'$

from x' **have** $y: y \in Y$ **unfolding** *Y-def* *y-def* **by** *auto*

note *id*

also **have** $r x' \leq r y$ **using** *r-pow-A-1*[*OF x'(2)*] **unfolding** *y-def* .

also **have** $\dots \leq r z$ **using** *sr-max-Y*[*OF y*] .

finally **show** $r x \leq r z$.

qed

lemma *z-pos*: $z \ \$ \ i > 0$

using *Y-pos*[*OF z(1)*] **by** *auto*

lemma *sr-pos*: $sr > 0$

by (*rule rx-pos, insert z-pos, auto*)

context **fixes** u

```

    assumes  $u: u \in X$  and  $ru: r u = sr$ 
  begin

  lemma sr-imp-eigen-vector-main:  $sr * s u = A * v u$ 
  proof (rule ccontr)
    assume *:  $sr * s u \neq A * v u$ 
    let  $?x = A * v u - sr * s u$ 
    from * have  $0: ?x \neq 0$  by auto
    let  $?y = pow-A-1 u$ 
    have le-vec ( $sr * s u$ ) ( $A * v u$ ) using rx-le-Ax[OF  $u$ ] unfolding ru .
    hence le: le-vec  $0$   $?x$  by auto
    from  $0$  le have  $x: ?x \in X$  unfolding X-def by auto
    have y-pos: lt-vec  $0$   $?y$  using Y-pos-main[of  $?y$ ]  $u$  by auto
    hence  $y: ?y \in X$  unfolding X-def by (auto simp: order.strict-iff-order)
    from Y-pos-main[of  $pow-A-1$   $?x$ ]  $x$ 
    have lt-vec  $0$  ( $pow-A-1$   $?x$ ) by auto
    hence lt: lt-vec ( $sr * s ?y$ ) ( $A * v ?y$ ) unfolding pow-A-1-def matrix-vector-right-distrib-diff
      matrix-vector-mul-assoc A1n-commute vector-smult-distrib by simp
    let  $?f = (\lambda i. (A * v ?y - sr * s ?y) \$ i / ?y \$ i)$ 
    let  $?U = UNIV :: 'n$  set
    define eps where  $eps = Min (?f ' ?U)$ 
    have  $U: finite (?f ' ?U)$   $?f ' ?U \neq \{\}$  by auto
    have eps:  $eps > 0$  unfolding eps-def Min-gr-iff[OF  $U$ ]
      using lt sr-pos y-pos by auto
    have le: le-vec ( $(sr + eps) * s ?y$ ) ( $A * v ?y$ )
  proof
    fix  $i$ 
    have  $((sr + eps) * s ?y) \$ i = sr * ?y \$ i + eps * ?y \$ i$ 
      by (simp add: comm-semiring-class.distrib)
    also have  $\dots \leq sr * ?y \$ i + ?f i * ?y \$ i$ 
    proof (rule add-left-mono[OF mult-right-mono])
      show  $0 \leq ?y \$ i$  using y-pos[rule-format, of  $i$ ] by auto
      show  $eps \leq ?f i$  unfolding eps-def by (rule Min-le, auto)
    qed
    also have  $\dots = (A * v ?y) \$ i$  using sr-pos y-pos[rule-format, of  $i$ ]
      by simp
    finally
      show  $((sr + eps) * s ?y) \$ i \leq (A * v ?y) \$ i$  .
  qed
  from rho-le-x-Ax-imp-rho-le-rx[OF  $y$  le]
  have  $r ?y \geq sr + eps$  .
  with sr-max[OF  $y$ ] eps show False by auto
  qed

  lemma sr-imp-eigen-vector: eigen-vector  $A$   $u$  sr
    unfolding eigen-vector-def sr-imp-eigen-vector-main using  $u$  unfolding X-def
  by auto

  lemma sr-u-pos: lt-vec  $0$   $u$ 

```



```

proof –
  let ?y = pow-A-1 u
  define n where n = N
  define c where c = (sr + 1) ^ N
  have c: c > 0 using sr-pos unfolding c-def by auto
  have lt-vec 0 ?y using Y-pos-main[of ?y] u by auto
  also have ?y = A1n *v u unfolding pow-A-1-def ..
  also have ... = c *s u unfolding c-def A1n-def n-def[symmetric]
  proof (induct n)
    case (Suc n)
    then show ?case
      by (simp add: matrix-vector-mul-assoc[symmetric] algebra-simps vec.scale
        sr-imp-eigen-vector-main[symmetric])
    qed auto
  finally have lt: lt-vec 0 (c *s u) .
  have 0 < u $ i for i using lt[rule-format, of i] c by simp (metis zero-less-mult-pos)
  thus lt-vec 0 u by simp
qed
end

```

```

lemma eigen-vector-z-sr: eigen-vector A z sr
  using sr-imp-eigen-vector[OF zX refl] by auto

```

```

lemma eigen-value-sr: eigen-value A sr
  using eigen-vector-z-sr unfolding eigen-value-def by auto

```

```

abbreviation c ≡ complex-of-real
abbreviation cA ≡ map-matrix c A
abbreviation norm-v ≡ map-vector (norm :: complex ⇒ real)

```

```

lemma norm-v-ge-0: le-vec 0 (norm-v v) by (auto simp: map-vector-def)
lemma norm-v-eq-0: norm-v v = 0 ⟷ v = 0 by (auto simp: map-vector-def
  vec-eq-iff)

```

```

lemma cA-index: cA $ i $ j = c (A $ i $ j)
  unfolding map-matrix-def map-vector-def by simp

```

```

lemma norm-cA[simp]: norm (cA $ i $ j) = A $ i $ j
  using nonneg[of i j] by (simp add: cA-index)

```

```

context fixes α v
  assumes ev: eigen-vector cA v α
begin

```

```

lemma evD: α *s v = cA *v v v ≠ 0
  using ev[unfolded eigen-vector-def] by auto

```

```

lemma ev-alpha-norm-v: norm-v (α *s v) = (norm α *s norm-v v)
  by (auto simp: map-vector-def norm-mult vec-eq-iff)

```

lemma *ev-A-norm-v*: $\text{norm-v } (cA *v v) \$ j \leq (A *v \text{norm-v } v) \$ j$
proof –
 have $\text{norm-v } (cA *v v) \$ j = \text{norm } (\sum_{i \in \text{UNIV}} cA \$ j \$ i * v \$ i)$
 unfolding *map-vector-def* **by** (*simp add: matrix-vector-mult-def*)
 also have $\dots \leq (\sum_{i \in \text{UNIV}} \text{norm } (cA \$ j \$ i * v \$ i))$ **by** (*rule norm-sum*)
 also have $\dots = (\sum_{i \in \text{UNIV}} A \$ j \$ i * \text{norm-v } v \$ i)$
 by (*rule sum.cong[OF refl], auto simp: norm-mult map-vector-def*)
 also have $\dots = (A *v \text{norm-v } v) \$ j$ **by** (*simp add: matrix-vector-mult-def*)
 finally show *?thesis* .
qed

lemma *ev-le-vec*: $\text{le-vec } (\text{norm } \alpha *s \text{norm-v } v) (A *v \text{norm-v } v)$
 using *arg-cong[OF evD(1), of norm-v, unfolded ev-alpha-norm-v]* *ev-A-norm-v*
by *auto*

lemma *norm-v-X*: $\text{norm-v } v \in X$
 using *norm-v-ge-0[of v] evD(2) norm-v-eq-0[of v]* **unfolding** *X-def* **by** *auto*

lemma *ev-inequalities*: $\text{norm } \alpha \leq r (\text{norm-v } v) \ r (\text{norm-v } v) \leq sr$
proof –
 have $v: \text{norm-v } v \in X$ **by** (*rule norm-v-X*)
 from *rho-le-x-Ax-imp-rho-le-rx[OF v ev-le-vec]*
 show $\text{norm } \alpha \leq r (\text{norm-v } v)$.
 from *sr-max[OF v]*
 show $r (\text{norm-v } v) \leq sr$.
qed

lemma *eigen-vector-norm-sr*: $\text{norm } \alpha \leq sr$ **using** *ev-inequalities* **by** *auto*
end

lemma *eigen-value-norm-sr*: **assumes** *eigen-value cA α*
 shows $\text{norm } \alpha \leq sr$
 using *eigen-vector-norm-sr[of - α] asms* **unfolding** *eigen-value-def* **by** *auto*

lemma *le-vec-trans*: $\text{le-vec } x \ y \implies \text{le-vec } y \ u \implies \text{le-vec } x \ u$
 using *order.trans[of x \$ i y \$ i u \$ i for i]* **by** *auto*

lemma *eigen-vector-z-sr-c*: *eigen-vector cA (map-vector c z) (c sr)*
 unfolding *of-real-hom.eigen-vector-hom* **by** (*rule eigen-vector-z-sr*)

lemma *eigen-value-sr-c*: *eigen-value cA (c sr)*
 using *eigen-vector-z-sr-c* **unfolding** *eigen-value-def* **by** *auto*

definition $w = \text{perron-frobenius.z } (\text{transpose } A)$

lemma *w*: $\text{transpose } A *v w = sr *s w \ \text{lt-vec } 0 \ w \ \text{perron-frobenius.sr } (\text{transpose } A) = sr$

proof –
interpret t : *perron-frobenius transpose* A
 by (rule *pf-transpose*)
from *eigen-vector-z-sr-c* t .*eigen-vector-z-sr-c*
have ev : *eigen-value* cA (c sr) *eigen-value* $t.cA$ (c $t.sr$)
 unfolding *eigen-value-def* **by** *auto*
 {
 fix x
have *eigen-value* $(t.cA)$ $x =$ *eigen-value* $(transpose\ cA)$ x
 unfolding *map-matrix-def* *map-vector-def* *transpose-def*
 by (*auto simp: vec-eq-iff*)
also have $\dots =$ *eigen-value* cA x **by** (rule *eigen-value-transpose*)
finally have *eigen-value* $(t.cA)$ $x =$ *eigen-value* cA x .
 } **note** $ev-id = this$
with ev **have** ev : *eigen-value* $t.cA$ (c sr) *eigen-value* cA (c $t.sr$) **by** *auto*
from *eigen-value-norm-sr*[*OF* $ev(2)$] *t.eigen-value-norm-sr*[*OF* $ev(1)$]
show id : $t.sr = sr$ **by** *auto*
from t .*eigen-vector-z-sr*[*unfolded id, folded w-def*] **show** $transpose\ A * v\ w = sr$
 $*s\ w$
 unfolding *eigen-vector-def* **by** *auto*
from t .*z-pos*[*folded w-def*] **show** $lt-vec\ 0\ w$ **by** *auto*
qed

lemma *c-mod-id*: $a \in \mathbb{R} \implies Re\ a \geq 0 \implies c\ (cmod\ a) = a$ **by** (*auto simp: Reals-def*)

lemma *pos-rowvector-mult-0*: **assumes** lt : $lt-vec\ 0\ x$
and 0 : ($rowvector\ x :: real^{n'} * v\ y = 0$ (**is** $?x * v - = 0$) **and** le : $le-vec\ 0\ y$
shows $y = 0$

proof –
 {
 fix i
assume $y\ \$\ i \neq 0$
with le **have** yi : $y\ \$\ i > 0$ **by** (*auto simp: order.strict-iff-order*)
have $0 = (?x * v\ y)\ \$\ i$ **unfolding** 0 **by** *simp*
also have $\dots = (\sum_{j \in UNIV} x\ \$\ j * y\ \$\ j)$
 unfolding *rowvector-def* *matrix-vector-mult-def* **by** *simp*
also have $\dots > 0$
 by (rule *sum-pos2*[*of* i], *insert yi lt le, auto intro!: mult-nonneg-nonneg simp: order.strict-iff-order*)
finally have *False* **by** *simp*
 }
thus $?thesis$ **by** (*auto simp: vec-eq-iff*)
qed

lemma *pos-matrix-mult-0*: **assumes** le : $\bigwedge i\ j. B\ \$\ i\ \$\ j \geq 0$
and lt : $lt-vec\ 0\ x$
and 0 : $B * v\ x = 0$

shows $B = 0$
proof –
{
 fix $i\ j$
 assume $B\ \$\ i\ \$\ j \neq 0$
 with le **have** $gt: B\ \$\ i\ \$\ j > 0$ **by** (*auto simp: order.strict-iff-order*)
 have $0 = (B *v x)\ \$\ i$ **unfolding** 0 **by** *simp*
 also have $\dots = (\sum_{j \in UNIV}. B\ \$\ i\ \$\ j * x\ \$\ j)$
 unfolding *matrix-vector-mult-def* **by** *simp*
 also have $\dots > 0$
 by (*rule sum-pos2[of - j], insert gt lt le, auto intro!: mult-nonneg-nonneg simp: order.strict-iff-order*)
 finally have *False* **by** *simp*
}
thus $B = 0$ **unfolding** *vec-eq-iff* **by** *auto*
qed

lemma *eigen-value-smaller-matrix*: **assumes** $B: \bigwedge i\ j. 0 \leq B\ \$\ i\ \$\ j \wedge B\ \$\ i\ \$\ j \leq A\ \$\ i\ \$\ j$
and $AB: A \neq B$
and ev : *eigen-value* (*map-matrix* $c\ B$) *sigma*
shows $cmod\ sigma < sr$
proof –
 let $?B = map-matrix\ c\ B$
 let $?sr = spectral-radius\ ?B$
 define σ **where** $\sigma = ?sr$
 have *real-non-neg-mat* $?B$ **unfolding** *real-non-neg-mat-def elements-mat-h-def*
 by (*auto simp: map-matrix-def map-vector-def B*)
 from *perron-frobenius[OF this, folded σ -def]* **obtain** x **where** $ev-sr$: *eigen-vector* $?B\ x\ (c\ \sigma)$
 and rnn : *real-non-neg-vec* x **by** *auto*
 define y **where** $y = norm-v\ x$
 from rnn **have** xy : $x = map-vector\ c\ y$
 unfolding *real-non-neg-vec-def vec-elements-h-def y-def*
 by (*auto simp: map-vector-def vec-eq-iff c-cmod-id*)
 from *spectral-radius-max[OF ev, folded σ -def]* **have** $sigma-sigma$: $cmod\ sigma \leq \sigma$.
 from $ev-sr$ [*unfolded xy of-real-hom.eigen-vector-hom*]
 have $ev-B$: *eigen-vector* $B\ y\ \sigma$.
 from $ev-B$ [*unfolded eigen-vector-def*] **have** $ev-B'$: $B *v\ y = \sigma *s\ y$ **by** *auto*
 have $ypos$: $y\ \$\ i \geq 0$ **for** i **unfolding** y -*def* **by** (*auto simp: map-vector-def*)
 from $ev-B$ *this* **have** y : $y \in X$ **unfolding** *eigen-vector-def X-def* **by** *auto*

 have BA : $(B *v\ y)\ \$\ i \leq (A *v\ y)\ \$\ i$ **for** i
 unfolding *matrix-vector-mult-def vec-lambda-beta*
 by (*rule sum-mono, rule mult-right-mono, insert B ypos, auto*)
 hence $le-vec$: $le-vec\ (\sigma *s\ y)\ (A *v\ y)$ **unfolding** $ev-B'$ **by** *auto*
 from *rho-le-x-Ax-imp-rho-le-rx[OF y le-vec]*
 have $\sigma \leq r\ y$ **by** *auto*

```

also have ... ≤ sr using y by (rule sr-max)
finally have sig-le-sr: σ ≤ sr .
{
  assume σ = sr
  hence r-sr: r y = sr and sr-sig: sr = σ using ⟨σ ≤ r y⟩ ⟨r y ≤ sr⟩ by auto
  from sr-u-pos[OF y r-sr] have pos: lt-vec 0 y .
  from sr-imp-eigen-vector[OF y r-sr] have ev': eigen-vector A y sr .
  have (A - B) *v y = A *v y - B *v y unfolding matrix-vector-mult-def
    by (auto simp: vec-eq-iff field-simps sum-subtractf)
  also have A *v y = sr *s y using ev'[unfolded eigen-vector-def] by auto
  also have B *v y = sr *s y unfolding ev-B' sr-sig ..
  finally have id: (A - B) *v y = 0 by simp
  from pos-matrix-mult-0[OF - pos id] assms(1-2) have False by auto
}
with sig-le-sr sigma-sigma show ?thesis by argo
qed

lemma charpoly-erase-mat-sr: 0 < poly (charpoly (erase-mat A i i)) sr
proof -
  let ?A = erase-mat A i i
  let ?pos = poly (charpoly ?A) sr
  {
    from A-nonzero-fixed-j[of i] obtain k where A $ k $ i ≠ 0 by auto
    assume A = ?A
    hence A $ k $ i = ?A $ k $ i by simp
    also have ?A $ k $ i = 0 by (auto simp: erase-mat-def)
    also have A $ k $ i ≠ 0 by fact
    finally have False by simp
  }
  hence AA: A ≠ ?A by auto
  have le: 0 ≤ ?A $ i $ j ∧ ?A $ i $ j ≤ A $ i $ j for i j
    by (auto simp: erase-mat-def nonneg)
  note ev-small = eigen-value-smaller-matrix[OF le AA]
  {
    fix rho :: real
    assume eigen-value ?A rho
    hence ev: eigen-value (map-matrix c ?A) (c rho)
      unfolding eigen-value-def using of-real-hom.eigen-vector-hom[of ?A - rho]
  }
by auto
  from ev-small[OF this] have abs rho < sr by auto
} note ev-small-real = this
have pos0: ?pos ≠ 0
  using ev-small-real[of sr] by (auto simp: eigen-value-root-charpoly)
{
  define p where p = charpoly ?A
  assume pos: ?pos < 0
  hence neg: poly p sr < 0 unfolding p-def by auto
  from degree-monic-charpoly[of ?A] have mon: monic p and deg: degree p ≠ 0
  unfolding p-def by auto
}

```

let $?f = \text{poly } p$
 have *cont*: *continuous-on* $\{a..b\}$ $?f$ for a b by (*auto intro*: *continuous-intros*)
 from *pos* have *le*: $?f \text{ sr} \leq 0$ by (*auto simp*: *p-def*)
 from *mon* have *lc*: *lead-coeff* $p > 0$ by *auto*
 from *poly-pinfty-ge*[*OF this deg, of 0*] obtain z where *lez*: $\bigwedge x. z \leq x \implies 0 \leq ?f x$ by *auto*
 define y where $y = \max z \text{ sr}$
 have *yr*: $y \geq \text{sr}$ and $y \geq z$ *unfolding* *y-def* by *auto*
 from *lez*[*OF this(2)*] have *y0*: $?f y \geq 0$.
 from *IVT*'[*of ?f, OF le y0 yr cont*] obtain x where *ge*: $x \geq \text{sr}$ and *rt*: $?f x = 0$
 unfolding *p-def* by *auto*
 hence *eigen-value* $?A x$ *unfolding* *p-def* by (*simp add*: *eigen-value-root-charpoly*)
 from *ev-small-real*[*OF this*] *ge* have *False* by *auto*
 }
 with *pos0* show *?thesis* by *argo*
 qed

lemma *multiplicity-sr-1*: *order sr (charpoly A) = 1*

proof –

{
 assume *poly (pderiv (charpoly A)) sr = 0*
 hence $0 = \text{poly} (\text{monom } 1 \ 1 * \text{pderiv} (\text{charpoly } A)) \text{ sr}$ by *simp*
 also have $\dots = \text{sum} (\lambda i. \text{poly} (\text{charpoly} (\text{erase-mat } A \ i \ i)) \text{ sr})$ *UNIV*
 unfolding *pderiv-char-poly-erase-mat poly-sum ..*
 also have $\dots > 0$
 by (*rule sum-pos, (force simp: charpoly-erase-mat-sr)+*)
 finally have *False* by *simp*
 }
 hence nZ : *poly (pderiv (charpoly A)) sr $\neq 0$* and nZ' : *pderiv (charpoly A) $\neq 0$*
 by *auto*
 from *eigen-vector-z-sr* have *eigen-value A sr* *unfolding* *eigen-value-def ..*
 from *this*[*unfolded eigen-value-root-charpoly*]
 have *poly (charpoly A) sr = 0* .
 hence *order sr (charpoly A) $\neq 0$* *unfolding* *order-root* using nZ' by *auto*
 from *order-pderiv*[*OF nZ' this*] *order-0I*[*OF nZ*]
 show *?thesis* by *simp*
 qed

lemma *sr-spectral-radius*: *sr = spectral-radius cA*

proof –

from *eigen-vector-z-sr-c* have *eigen-value cA (c sr)*
 unfolding *eigen-value-def* by *auto*
 from *spectral-radius-max*[*OF this*]
 have *sr*: $\text{sr} \leq \text{spectral-radius } cA$ by *auto*
 with *spectral-radius-ev*[*of cA*] *eigen-vector-norm-sr*
 show *?thesis* by *force*
 qed

lemma *le-vec-A-mu*: **assumes** $y: y \in X$ **and** $le: le\text{-vec } (A * v y) (mu * s y)$
shows $sr \leq mu$ *lt-vec 0 y*
 $mu = sr \vee A * v y = mu * s y \implies mu = sr \wedge A * v y = mu * s y$

proof –

let $?w = \text{rowvector } w$
let $?w' = \text{columnvector } w$
have $?w ** A = \text{transpose } (\text{transpose } (?w ** A))$
unfolding *transpose-transpose* **by** *simp*
also have $\text{transpose } (?w ** A) = \text{transpose } A ** \text{transpose } ?w$
by *(rule matrix-transpose-mul)*
also have $\text{transpose } ?w = \text{columnvector } w$ **by** *(rule transpose-rowvector)*
also have $\text{transpose } A ** \dots = \text{columnvector } (\text{transpose } A * v w)$
unfolding *dot-rowvector-columnvector[symmetric]* **..**
also have $\text{transpose } A * v w = sr * s w$ **unfolding** w **by** *simp*
also have $\text{transpose } (\text{columnvector } \dots) = \text{rowvector } (sr * s w)$
unfolding *transpose-def columnvector-def rowvector-def vector-scalar-mult-def*
by *auto*
finally have $1: ?w ** A = \text{rowvector } (sr * s w)$.
have $sr * s (?w * v y) = ?w ** A * v y$ **unfolding** 1
by *(auto simp: rowvector-def vector-scalar-mult-def matrix-vector-mult-def vec-eq-iff sum-distrib-left mult.assoc)*
also have $\dots = ?w * v (A * v y)$ **by** *(simp add: matrix-vector-mult-assoc)*
finally have $eq1: sr * s (\text{rowvector } w * v y) = \text{rowvector } w * v (A * v y)$.
have $le\text{-vec } (\text{rowvector } w * v (A * v y)) (?w * v (mu * s y))$
by *(rule le-vec-mono-left[OF - le], insert w(2), auto simp: rowvector-def order.strict-iff-order)*
also have $?w * v (mu * s y) = mu * s (?w * v y)$ **by** *(simp add: algebra-simps vec.scale)*
finally have $le1: le\text{-vec } (\text{rowvector } w * v (A * v y)) (mu * s (?w * v y))$.
from $le1$ *[unfolded eq1[symmetric]]*
have $2: le\text{-vec } (sr * s (?w * v y)) (mu * s (?w * v y))$.
{
from y **obtain** i **where** $y_i: y \$ i > 0$ **and** $y: \bigwedge j. y \$ j \geq 0$ **unfolding** $X\text{-def}$
by *(auto simp: order.strict-iff-order vec-eq-iff)*
from $w(2)$ **have** $w_i: w \$ i > 0$ **and** $w: \bigwedge j. w \$ j \geq 0$
by *(auto simp: order.strict-iff-order)*
have $(?w * v y) \$ i > 0$ **using** $y_i y w$
by *(auto simp: matrix-vector-mult-def rowvector-def intro!: sum-pos2[of - i] mult-nonneg-nonneg)*
moreover from 2 *[rule-format, of i]* **have** $sr * (?w * v y) \$ i \leq mu * (?w * v y) \$ i$ **by** *simp*
ultimately have $sr \leq mu$ **by** *simp*
}
thus $*$: $sr \leq mu$.
define cc **where** $cc = (mu + 1) ^ N$
define n **where** $n = N$
from $*$ *sr-pos* **have** $mu: mu \geq 0$ $mu > 0$ **by** *auto*
hence $cc: cc > 0$ **unfolding** $cc\text{-def}$ **by** *simp*
from y **have** $\text{pow-}A\text{-}1 y \in \text{pow-}A\text{-}1 ' X$ **by** *auto*

from *Y-pos-main*[*OF this*] **have** $lt: 0 < (A1n * v y) \$ i$ **for** i **by** (*simp add: pow-A-1-def*)
have $le: le-vec (A1n * v y) (cc * s y)$ **unfolding** *cc-def A1n-def n-def*[*symmetric*]
proof (*induct n*)
case (*Suc n*)
let $?An = matpow (A + mat 1) n$
let $?mu = (mu + 1)$
have $id': matpow (A + mat 1) (Suc n) * v y = A * v (?An * v y) + ?An * v y$
(is $?a = ?b + ?c$)
by (*simp add: matrix-add-ldistrib matrix-mul-rid matrix-add-vect-distrib mat-pow-1-commute*
matrix-vector-mul-assoc[*symmetric*])
have $le-vec ?b (?mu \hat{n} * s (A * v y))$
using *le-vec-mono-left*[*OF nonneg Suc*] **by** (*simp add: algebra-simps vec.scale*)
moreover **have** $le-vec (?mu \hat{n} * s (A * v y)) (?mu \hat{n} * s (mu * s y))$
using *le mu by auto*
moreover **have** $id: ?mu \hat{n} * s (mu * s y) = (?mu \hat{n} * mu) * s y$ **by** *simp*
from *le-vec-trans*[*OF calculation*[*unfolded id*]]
have $le1: le-vec ?b ((?mu \hat{n} * mu) * s y)$.
from *Suc* **have** $le2: le-vec ?c ((mu + 1) \hat{n} * s y)$.
have $le: le-vec ?a ((?mu \hat{n} * mu) * s y + ?mu \hat{n} * s y)$
unfolding id' **using** *add-mono*[*OF le1*[*rule-format*] *le2*[*rule-format*]] **by** *auto*
have $id'': (?mu \hat{n} * mu) * s y + ?mu \hat{n} * s y = ?mu \hat{Suc} n * s y$ **by** (*simp add: algebra-simps*)
show $?case$ **using** *le* **unfolding** id'' .
qed (*simp add: matrix-vector-mul-lid*)
have $lt: 0 < cc * y \$ i$ **for** i **using** $lt[of i]$ le [*rule-format, of i*] **by** *auto*
have $y \$ i > 0$ **for** i **using** $lt[of i]$ cc **by** (*rule zero-less-mult-pos*)
thus $lt-vec 0 y$ **by** *auto*
assume $**$: $mu = sr \vee A * v y = mu * s y$
{
assume $A * v y = mu * s y$
with y **have** *eigen-vector* $A y mu$ **unfolding** *X-def eigen-vector-def* **by** *auto*
hence *eigen-vector* $cA (map-vector c y) (c mu)$ **unfolding** *of-real-hom.eigen-vector-hom*
.

from *eigen-vector-norm-sr*[*OF this*] **have** $mu = sr$ **by** *auto*
}
with $**$ **have** $mu-sr: mu = sr$ **by** *auto*
from *eq1*[*folded vector-smult-distrib*]
have $0: ?w * v (sr * s y - A * v y) = 0$
unfolding *matrix-vector-right-distrib-diff* **by** *simp*
have $le0: le-vec 0 (sr * s y - A * v y)$ **using** *assms(2)*[*unfolded mu-sr*] **by** *auto*
have $sr * s y - A * v y = 0$ **using** *pos-rowvector-mult-0*[*OF w(2) 0 le0*] .
hence $ev-y: A * v y = sr * s y$ **by** *auto*
show $mu = sr \wedge A * v y = mu * s y$ **using** $ev-y mu-sr$ **by** *auto*
qed

lemma *nonnegative-eigenvector-has-ev-sr*: **assumes** *eigen-vector* $A v mu$ **and** $le: le-vec 0 v$


```

shows mu = sr
proof -
  from assms(1)[unfolded eigen-vector-def] have v: v ≠ 0 and ev: A *v v = mu
  *s v by auto
  from le v have v: v ∈ X unfolding X-def by auto
  from ev have le-vec (A *v v) (mu *s v) by auto
  from le-vec-A-mu[OF v this] ev show ?thesis by auto
qed

lemma similar-matrix-rotation: assumes ev: eigen-value cA α and α: cmod α =
sr
shows similar-matrix (cis (Arg α) *k cA) cA
proof -
  from ev obtain y where ev: eigen-vector cA y α unfolding eigen-value-def by
  auto
  let ?y = norm-v y
  note maps = map-vector-def map-matrix-def
  define yp where yp = norm-v y
  let ?yp = map-vector c yp
  have yp: yp ∈ X unfolding yp-def by (rule norm-v-X[OF ev])
  from ev[unfolded eigen-vector-def] have ev-y: cA *v y = α *s y by auto
  from ev-le-vec[OF ev, unfolded α, folded yp-def]
  have 1: le-vec (sr *s yp) (A *v yp) by simp
  from rho-le-x-Ax-imp-rho-le-rx[OF yp 1] have sr ≤ r yp by auto
  with ev-inequalities[OF ev, folded yp-def]
  have 2: r yp = sr by auto
  have ev-yp: A *v yp = sr *s yp
    and pos-yp: lt-vec 0 yp
    using sr-imp-eigen-vector-main[OF yp 2] sr-u-pos[OF yp 2] by auto
  define D where D = diagvector (λ j. cis (Arg (y $ j)))
  define inv-D where inv-D = diagvector (λ j. cis (- Arg (y $ j)))
  have DD: inv-D ** D = mat 1 D ** inv-D = mat 1 unfolding D-def inv-D-def
  by (auto simp add: diagvector-eq-mat cis-mult)
  {
    fix i
    have (D *v ?yp) $ i = cis (Arg (y $ i)) * c (cmod (y $ i))
      unfolding D-def yp-def by (simp add: maps)
    also have ... = y $ i by (simp add: cis-mult-cmod-id)
    also note calculation
  }
  hence y-D-yp: y = D *v ?yp by (auto simp: vec-eq-iff)
  define φ where φ = Arg α
  let ?φ = cis (- φ)
  have [simp]: cis (- φ) * rcis sr φ = sr unfolding cis-rcis-eq rcis-mult by simp
  have α: α = rcis sr φ unfolding φ-def α[symmetric] rcis-cmod-Arg ..
  define F where F = ?φ *k (inv-D ** cA ** D)
  have cA *v (D *v ?yp) = α *s y unfolding y-D-yp[symmetric] ev-y by simp
  also have inv-D *v ... = α *s ?yp
    unfolding vector-smult-distrib y-D-yp matrix-vector-mul-assoc DD matrix-vector-mul-lid

```

..

also have $?φ * s \dots = sr * s ?yp$ **unfolding** $α$ **by** *simp*

also have $\dots = \text{map-vector } c (sr * s yp)$ **unfolding** *vec-eq-iff* **by** (*auto simp: maps*
maps)

also have $\dots = cA * v ?yp$ **unfolding** *ev-yp[symmetric]* **by** (*auto simp: maps*
matrix-vector-mult-def)

finally have $F: F * v ?yp = cA * v ?yp$ **unfolding** *F-def matrix-scalar-vector-ac[symmetric]*
unfolding matrix-vector-mul-assoc[symmetric] vector-smult-distrib .

have *prod: inv-D ** cA ** D = (χ i j. cis (- Arg (y \$ i)) * cA \$ i \$ j * cis*
(Arg (y \$ j)))

unfolding *inv-D-def D-def diagvector-mult-right diagvector-mult-left* **by** *simp*

{

 fix *i j*

 have *cmod (F \$ i \$ j) = cmod (?φ * cA \$h i \$h j * (cis (- Arg (y \$h i)) * cis*
(Arg (y \$h j))))

unfolding *F-def prod vec-lambda-beta matrix-scalar-mult-def*

by (*simp only: ac-simps*)

 also have $\dots = A \$ i \$ j$ **unfolding** *cis-mult* **unfolding** *norm-mult* **by** *simp*

 also **note** *calculation*

}

hence *FA: map-matrix norm F = A* **unfolding** *maps* **by** *auto*

let $?F = \text{map-matrix } c (\text{map-matrix norm } F)$

let $?G = ?F - F$

let $?Re = \text{map-matrix } Re$

from *F[folded FA]* **have** $0: ?G * v ?yp = 0$ **unfolding** *matrix-diff-vect-distrib* **by**
simp

have $?Re ?G * v yp = \text{map-vector } Re (?G * v ?yp)$

unfolding *maps matrix-vector-mult-def vec-lambda-beta Re-sum* **by** *auto*

also have $\dots = 0$ **unfolding** *0* **by** (*simp add: vec-eq-iff maps*)

finally have $0: ?Re ?G * v yp = 0$.

have $?Re ?G = 0$

by (*rule pos-matrix-mult-0[OF - pos-yp 0], auto simp: maps complex-Re-le-cmod*)

hence $?F = F$ **by** (*auto simp: maps vec-eq-iff cmod-eq-Re*)

with *FA* **have** $AF: cA = F$ **by** *simp*

from *arg-cong[OF this, of λ A. cis φ *k A]*

have *sim: cis φ *k cA = inv-D ** cA ** D* **unfolding** *F-def matrix.scale-scale*
cis-mult

by *simp*

have *similar-matrix (cis φ *k cA) cA* **unfolding** *similar-matrix-def similar-matrix-wit-def*
sim

by (*rule exI[of - inv-D], rule exI[of - D], auto simp: DD*)

thus *?thesis* **unfolding** *φ-def* .

qed

lemma assumes *ev: eigen-value cA α* **and** $α: \text{cmod } α = sr$

shows *maximal-eigen-value-order-1: order α (charpoly cA) = 1*

and *maximal-eigen-value-rotation: eigen-value cA (x * cis (Arg α)) = eigen-value*
cA x

eigen-value cA (x / cis (Arg α)) = eigen-value cA x

proof –

let $?a = \text{cis } (\text{Arg } \alpha)$
let $?p = \text{charpoly } cA$
from *similar-matrix-rotation*[*OF ev α*]
have *similar-matrix* ($?a *k cA$) cA .
from *similar-matrix-charpoly*[*OF this*]
have *id*: *charpoly* ($?a *k cA$) = $?p$.
have $a: ?a \neq 0$ **by** *simp*
from *order-charpoly-smult*[*OF this, of - cA, unfolded id*]
have *order-neg*: *order* $x ?p = \text{order } (x / ?a) ?p$ **for** x .
have *order-pos*: *order* $x ?p = \text{order } (x * ?a) ?p$ **for** x
using *order-neg*[*symmetric, of $x * ?a$*] **by** *simp*
note *order-neg*[*of α*]
also **have** *id*: $\alpha / ?a = \text{sr}$ **unfolding** α [*symmetric*]
by (*metis a cis-mult-cmod-id nonzero-mult-div-cancel-left*)
also **have** *sr*: *order* ... $?p = 1$ **unfolding** *multiplicity-sr-1*[*symmetric*] *char-*
poly-of-real
by (*rule map-poly-inj-idom-divide-hom.order-hom, unfold-locales*)
finally **show** $*$: *order* $\alpha ?p = 1$.
show *eigen-value* $cA (x * ?a) = \text{eigen-value } cA x$ **using** *order-pos*
unfolding *eigen-value-root-charpoly order-root* **by** *auto*
show *eigen-value* $cA (x / ?a) = \text{eigen-value } cA x$ **using** *order-neg*
unfolding *eigen-value-root-charpoly order-root* **by** *auto*
qed

lemma *maximal-eigen-values-group*: **assumes** $M: M = \{ev :: \text{complex. eigen-value } cA \text{ ev} \wedge \text{cmod } ev = \text{sr}\}$
and $a: \text{rcis } sr \alpha \in M$
and $b: \text{rcis } sr \beta \in M$
shows $\text{rcis } sr (\alpha + \beta) \in M$ $\text{rcis } sr (\alpha - \beta) \in M$ $\text{rcis } sr 0 \in M$
proof –
{
fix a
assume $*$: $\text{rcis } sr a \in M$
have *id*: $\text{cis } (\text{Arg } (\text{rcis } sr a)) = \text{cis } a$
by (*smt * M mem-Collect-eq nonzero-mult-div-cancel-left of-real-eq-0-iff*
rcis-cmod-Arg rcis-def sr-pos)
from $*$ [*unfolded assms*] **have** *eigen-value* $cA (\text{rcis } sr a) \text{cmod } (\text{rcis } sr a) = \text{sr}$
by *auto*
from *maximal-eigen-value-rotation*[*OF this, unfolded id*]
have *eigen-value* $cA (x * \text{cis } a) = \text{eigen-value } cA x$
eigen-value $cA (x / \text{cis } a) = \text{eigen-value } cA x$ **for** x **by** *auto*
} **note** $*$ = *this*
from $*$ (1)[*OF b, of rcis sr α*] a **show** $\text{rcis } sr (\alpha + \beta) \in M$ **unfolding** M **by**
auto
from $*$ (2)[*OF a, of rcis sr α*] a **show** $\text{rcis } sr 0 \in M$ **unfolding** M **by** *auto*
from $*$ (2)[*OF b, of rcis sr α*] a **show** $\text{rcis } sr (\alpha - \beta) \in M$ **unfolding** M **by**
auto
qed

lemma *maximal-eigen-value-roots-of-unity-rotation*:

assumes $M: M = \{ev :: \text{complex. eigen-value } cA \text{ ev} \wedge \text{cmod } ev = sr\}$
and $kM: k = \text{card } M$

shows $k \neq 0$
 $k \leq \text{CARD}('n)$
 $\exists f. \text{charpoly } A = (\text{monom } 1 \ k - [:sr \wedge k:]) * f$
 $\wedge (\forall x. \text{poly } (\text{map-poly } c \ f) \ x = 0 \longrightarrow \text{cmod } x < sr)$
 $M = (*) (c \ sr) \ ' (\lambda i. (\text{cis } (\text{of-nat } i * 2 * \pi / k))) \ ' \{0 ..< k\}$
 $M = (*) (c \ sr) \ ' \{x :: \text{complex. } x \wedge k = 1\}$
 $(*) (\text{cis } (2 * \pi / k)) \ ' \text{Spectrum } cA = \text{Spectrum } cA$

unfolding kM

proof –

let $?M = \text{card } M$
note $fin = \text{finite-spectrum}[of \ cA]$
note $char = \text{degree-monic-charpoly}[of \ cA]$
have $?M \leq \text{card } (\text{Collect } (\text{eigen-value } cA))$
by $(\text{rule } \text{card-mono}[OF \ fin], \text{unfold } M, \text{auto})$
also have $\text{Collect } (\text{eigen-value } cA) = \{x. \text{poly } (\text{charpoly } cA) \ x = 0\}$
unfolding *eigen-value-root-charpoly* **by** *auto*
also have $\text{card } \dots \leq \text{degree } (\text{charpoly } cA)$
by $(\text{rule } \text{poly-roots-degree}, \text{insert } char, \text{auto})$
also have $\dots = \text{CARD}('n)$ **using** *char* **by** *simp*
finally show $?M \leq \text{CARD}('n)$.
from *finite-subset*[*OF* - *fin*, *of* *M*]
have $finM: \text{finite } M$ **unfolding** *M* **by** *blast*
from *finite-distinct-list*[*OF* *this*]
obtain m **where** $Mm: M = \text{set } m$ **and** $dist: \text{distinct } m$ **by** *auto*
from $Mm \ dist$ **have** $card: ?M = \text{length } m$ **by** $(\text{auto } \text{simp}: \text{distinct-card})$
have $sr: sr \in \text{set } m$ **using** *eigen-value-sr-c sr-pos* **unfolding** $Mm[\text{symmetric}] \ M$
by *auto*
define s **where** $s = \text{sort-key } Arg \ m$
define a **where** $a = \text{map } Arg \ s$
let $?k = \text{length } a$
from $dist \ Mm \ \text{card } sr$ **have** $s: M = \text{set } s$ $\text{distinct } s$ $sr \in \text{set } s$
and $card: ?M = ?k$
and $sorted: \text{sorted } a$
unfolding *s-def a-def* **by** *auto*
have $map-s: \text{map } ((*) (c \ sr)) (\text{map } cis \ a) = s$ **unfolding** *map-map o-def a-def*
proof $(\text{rule } \text{map-idI})$
fix x
assume $x \in \text{set } s$
from $this[\text{folded } s(1), \text{unfolded } M]$
have $id: \text{cmod } x = sr$ **by** *auto*
show $sr * cis (Arg \ x) = x$
by $(\text{subst } (5) \ \text{rcis-cmod-Arg}[\text{symmetric}], \text{unfold } id[\text{symmetric}] \ \text{rcis-def}, \text{simp})$

qed

from $s(2)[\text{folded } map-s, \text{unfolded } \text{distinct-map}]$ **have** $a: \text{distinct } a$ *inj-on* *cis* $(\text{set } a)$ **by** *auto*

from $s(3)$ **obtain** $aa\ a'$ **where** a -split: $a = aa \# a'$ **unfolding** a -def **by** (*cases*
s, auto)
from *Arg-bounded* **have** $\text{bounded}: x \in \text{set } a \implies -\text{pi} < x \wedge x \leq \text{pi}$ **for** x
unfolding a -def **by** *auto*
from $\text{bounded}[\text{of } aa, \text{unfolded } a\text{-split}]$ **have** $aa: -\text{pi} < aa \wedge aa \leq \text{pi}$ **by** *auto*
let $?aa = aa + 2 * \text{pi}$
define args **where** $\text{args} = a @ [?aa]$
let $?diff = \lambda i. \text{args} ! \text{Suc } i - \text{args} ! i$
have $\text{bnd}: x \in \text{set } a \implies x < ?aa$ **for** x **using** aa $\text{bounded}[\text{of } x]$ **by** *auto*
hence aa - a : $?aa \notin \text{set } a$ **by** *fast*
have $\text{sorted}: \text{sorted } \text{args}$ **unfolding** args-def **using** *sorted* **unfolding** *sorted-append*
by (*insert bnd, auto simp: order.strict-iff-order*)
have $\text{dist}: \text{distinct } \text{args}$ **using** aa - a **unfolding** args-def *distinct-append* **by** *auto*
have $\text{sum}: (\sum i < ?k. ?diff\ i) = 2 * \text{pi}$
unfolding *sum-lessThan-telescope* args-def a -split **by** *simp*
have $k: ?k \neq 0$ **unfolding** a -split **by** *auto*
let $?A = ?diff\ ' \{.. < ?k\}$
let $?Min = \text{Min } ?A$
define Min **where** $\text{Min} = ?Min$
have $?Min = (?k * ?Min) / ?k$ **using** k **by** *auto*
also **have** $?k * ?Min = (\sum i < ?k. ?Min)$ **by** *auto*
also **have** $\dots / ?k \leq (\sum i < ?k. ?diff\ i) / ?k$
by (*rule divide-right-mono[OF sum-mono[OF Min-le]]*, *auto*)
also **have** $\dots = 2 * \text{pi} / ?k$ **unfolding** sum **..**
finally **have** $\text{Min}: \text{Min} \leq 2 * \text{pi} / ?k$ **unfolding** Min-def **by** *auto*
have $\text{lt}: i < ?k \implies \text{args} ! i < \text{args} ! (\text{Suc } i)$ **for** i
using *sorted[unfolded sorted-iff-nth-mono, rule-format, of i Suc i]*
dist[unfolded distinct-conv-nth, rule-format, of Suc i i] **by** (*auto simp: args-def*)
let $?c = \lambda i. \text{rcis } \text{sr } (\text{args} ! i)$
have $\text{hda}[\text{simp}]: \text{hd } a = aa$ **unfolding** a -split **by** *simp*
have $\text{Min}0: \text{Min} > 0$ **using** lt **unfolding** Min-def **by** (*subst Min-gr-iff, insert*
k, auto)
have $\text{Min-A}: \text{Min} \in ?A$ **unfolding** Min-def **by** (*rule Min-in, insert k, auto*)
{
fix $i :: \text{nat}$
assume $i: i < \text{length } \text{args}$
hence $?c\ i = \text{rcis } \text{sr } ((a @ [\text{hd } a]) ! i)$
by (*cases i = ?k, auto simp: args-def nth-append rcis-def*)
also **have** $\dots \in \text{set } (\text{map } (\text{rcis } \text{sr}) (a @ [\text{hd } a]))$ **using** i
unfolding args-def *set-map* **unfolding** *set-conv-nth* **by** *auto*
also **have** $\dots = \text{rcis } \text{sr}\ ' \text{set } a$ **unfolding** a -split **by** *auto*
also **have** $\dots = M$ **unfolding** $s(1)$ *map-s[symmetric]* *set-map image-image*
by (*rule image-cong[OF refl], auto simp: rcis-def*)
finally **have** $?c\ i \in M$ **by** *auto*
} **note** $\text{ci}M = \text{this}$
{
fix $i :: \text{nat}$
assume $i: i < ?k$
hence $i < \text{length } \text{args}$ $\text{Suc } i < \text{length } \text{args}$ **unfolding** args-def **by** *auto*

```

    from maximal-eigen-values-group[OF M ciM[OF this(2)]] ciM[OF this(1)]
    have rcis sr (?diff i) ∈ M by simp
  }
  hence Min-M: rcis sr Min ∈ M using Min-A by force
  have rcisM: rcis sr (of-nat n * Min) ∈ M for n
  proof (induct n)
    case 0
    show ?case using sr Mm by auto
  next
    case (Suc n)
    have *: rcis sr (of-nat (Suc n) * Min) = rcis sr (of-nat n * Min) * cis Min
      by (simp add: rcis-mult ring-distrib add.commute)
    from maximal-eigen-values-group(1)[OF M Suc Min-M]
    show ?case unfolding * by simp
  qed
  let ?list = map (rcis sr) (map (λ i. of-nat i * Min) [0 ..< ?k])
  define list where list = ?list
  have len: length ?list = ?M unfolding card by simp
  from sr-pos have sr0: sr ≠ 0 by auto
  {
    fix i
    assume i: i < ?k
    hence *: 0 ≤ real i * Min using Min0 by auto
    from i have real i < real ?k by auto
    from mult-strict-right-mono[OF this Min0]
    have real i * Min < real ?k * Min by simp
    also have ... ≤ real ?k * (2 * pi / real ?k)
      by (rule mult-left-mono[OF Min], auto)
    also have ... = 2 * pi using k by simp
    finally have real i * Min < 2 * pi .
    note * this
  }
  note prod-pi = this
  have dist: distinct ?list
    unfolding distinct-map[of rcis sr]
  proof (rule conjI[OF - inj-on-subset[OF rcis-inj-on[OF sr0]]])
    show distinct (map (λ i. of-nat i * Min) [0 ..< ?k]) using Min0
      by (auto simp: distinct-map inj-on-def)
    show set (map (λ i. real i * Min) [0..<?k]) ⊆ {0..<2 * pi} using prod-pi
      by auto
  qed
  with len have card!: card (set ?list) = ?M using distinct-card by fastforce
  have listM: set ?list ⊆ M using rcisM by auto
  from card-subset-eq[OF finM listM card!]
  have M-list: M = set ?list ..
  let ?piM = 2 * pi / ?M
  {
    assume Min ≠ ?piM
    with Min have lt: Min < 2 * pi / ?k unfolding card by simp
    from k have 0 < real ?k by auto
  }

```

```

from mult-strict-left-mono[OF lt this] k Min0
have k: 0 ≤ ?k * Min ?k * Min < 2 * pi by auto
from rcisM[of ?k, unfolded M-list] have rcis sr (?k * Min) ∈ set ?list by auto
then obtain i where i: i < ?k and id: rcis sr (?k * Min) = rcis sr (i * Min)
by auto
from inj-onD[OF inj-on-subset[OF rcis-inj-on[OF sr0], of {?k * Min, i * Min}]]
id]
  prod-pi[OF i] k
  have ?k * Min = i * Min by auto
  with Min0 i have False by auto
}
hence Min: Min = ?piM by auto
show cM: ?M ≠ 0 unfolding card using k by auto
let ?f = (λ i. cis (of-nat i * 2 * pi / ?M))
note M-list
also have set ?list = (*) (c sr) ‘ (λ i. cis (of-nat i * Min)) ‘ {0 ..< ?k}
  unfolding set-map image-image
  by (rule image-cong, insert sr-pos, auto simp: rcis-mult rcis-def)
finally show M-cis: M = (*) (c sr) ‘ ?f ‘ {0 ..< ?M}
  unfolding card Min by (simp add: mult.assoc)
thus M-pow: M = (*) (c sr) ‘ {x :: complex. x ^ ?M = 1} using roots-of-unity[OF
cM] by simp
let ?rphi = rcis sr (2 * pi / ?M)
let ?phi = cis (2 * pi / ?M)
from Min-M[unfolded Min]
have ev: eigen-value cA ?rphi unfolding M by auto
have cm: cmod ?rphi = sr using sr-pos by simp
have id: cis (Arg ?rphi) = cis (Arg ?phi) * cmod ?phi
  unfolding arg-rcis-cis[OF sr-pos] by simp
also have ... = ?phi unfolding cis-mult-cmod-id ..
finally have id: cis (Arg ?rphi) = ?phi .
define phi where phi = ?phi
have phi: phi ≠ 0 unfolding phi-def by auto
note max = maximal-eigen-value-rotation[OF ev cm, unfolded id phi-def[symmetric]]
have ((* phi) ‘ Spectrum cA = Spectrum cA (is ?L = ?R)
proof –
  {
    fix x
    have *: x ∈ ?L ⇒ x ∈ ?R for x using max(2)[of x] phi unfolding
Spectrum-def by auto
    moreover
    {
      assume x ∈ ?R
      hence eigen-value cA x unfolding Spectrum-def by auto
      from this[folded max(2)[of x]] have x / phi ∈ ?R unfolding Spectrum-def
by auto
      from imageI[OF this, of (*) phi]
      have x ∈ ?L using phi by auto
    }
  }

```

```

    note this *
  }
  thus ?thesis by blast
qed
from this[unfolded phi-def]
show (*) (cis (2 * pi / real (card M))) ‘ Spectrum cA = Spectrum cA .
let ?p = monom 1 k - [:sr ^k:]
let ?cp = monom 1 k - [(c sr) ^k:]
let ?one = 1 :: complex
let ?list = map (rcis sr) (map (λ i. of-nat i * ?piM) [0 ..< card M])
interpret c: field-hom c ..
interpret p: map-poly-inj-idom-divide-hom c ..
have cp: ?cp = map-poly c ?p by (simp add: hom-distrib)
have M-list: M = set ?list using M-list[unfolded Min card[symmetric]] .
have dist: distinct ?list using dist[unfolded Min card[symmetric]] .
have k0: k ≠ 0 using k[folded card] assms by auto
have ?cp = (monom 1 k + (- [:(c sr) ^k:])) by simp
also have degree ... = k
  by (subst degree-add-eq-left, insert k0, auto simp: degree-monom-eq)
finally have deg: degree ?cp = k .
from deg k0 have cp0: ?cp ≠ 0 by auto
have {x. poly ?cp x = 0} = {x. x ^k = (c sr) ^k} unfolding poly-diff poly-monom

  by simp
also have ... ⊆ M
proof -
  {
    fix x
    assume id: x ^k = (c sr) ^k
    from sr-pos k0 have (c sr) ^k ≠ 0 by auto
    with arg-cong[OF id, of λ x. x / (c sr) ^k]
    have (x / c sr) ^k = 1
      unfolding power-divide by auto
    hence c sr * (x / c sr) ∈ M
      by (subst M-pow, unfold kM[symmetric], blast)
    also have c sr * (x / c sr) = x using sr-pos by auto
    finally have x ∈ M .
  }
  thus ?thesis by auto
qed
finally have cp-M: {x. poly ?cp x = 0} ⊆ M .
have k = card (set ?list) unfolding distinct-card[OF dist] by (simp add: kM)
also have ... ≤ card {x. poly ?cp x = 0}
proof (rule card-mono[OF poly-roots-finite[OF cp0]])
  {
    fix x
    assume x ∈ set ?list
    then obtain i where x: x = rcis sr (real i * ?piM) by auto
    have x ^k = (c sr) ^k unfolding x DeMoivre2 kM
  }

```



```

    by simp (metis mult.assoc of-real-power rcis-times-2pi)
  hence poly ?cp x = 0 unfolding poly-diff poly-monom by simp
}
thus set ?list ⊆ {x. poly ?cp x = 0} by auto
qed
finally have k-card: k ≤ card {x. poly ?cp x = 0} .
from k-card cp-M finM have M-id: M = {x. poly ?cp x = 0}
  unfolding kM by (metis card-seteq)
have dvdC: ?cp dvd charpoly cA
proof (rule poly-roots-dvd[OF cp0 deg k-card])
  from cp-M
  show {x. poly ?cp x = 0} ⊆ {x. poly (charpoly cA) x = 0}
    unfolding M eigen-value-root-charpoly by auto
qed
from this[unfolded charpoly-of-real cp p.hom-dvd-iff]
have dvd: ?p dvd charpoly A .
from this[unfolded dvd-def] obtain f where
  decomp: charpoly A = ?p * f by blast
let ?f = map-poly c f
have decompC: charpoly cA = ?cp * ?f unfolding charpoly-of-real decomp p.hom-mult
cp ..
show ∃ f. charpoly A = (monom 1 ?M - [:sr^?M:]) * f ∧ (∀ x. poly (map-poly
c f) x = 0 → cmod x < sr)
  unfolding kM[symmetric]
proof (intro exI conjI allI impI, rule decomp)
  fix x
  assume f: poly ?f x = 0
  hence ev: eigen-value cA x
    unfolding decompC p.hom-mult eigen-value-root-charpoly by auto
  hence le: cmod x ≤ sr using eigen-value-norm-sr by auto
  {
    assume max: cmod x = sr
    hence x ∈ M unfolding M using ev by auto
    hence poly ?cp x = 0 unfolding M-id by auto
    hence dvd1: [: -x, 1 :] dvd ?cp unfolding poly-eq-0-iff-dvd by auto
    from f[unfolded poly-eq-0-iff-dvd]
    have dvd2: [: -x, 1 :] dvd ?f by auto
    from char have 0: charpoly cA ≠ 0 by auto
    from mult-dvd-mono[OF dvd1 dvd2] have [: -x, 1 :]^2 dvd (charpoly cA)
      unfolding decompC power2-eq-square .
    from order-max[OF this 0] maximal-eigen-value-order-1[OF ev max]
    have False by auto
  }
  with le show cmod x < sr by argo
qed
qed
lemmas pf-main =
  eigen-value-sr eigen-vector-z-sr

```

```

eigen-value-norm-sr
z-pos
multiplicity-sr-1
nonnegative-eigenvector-has-ev-sr
maximal-eigen-value-order-1
maximal-eigen-value-roots-of-unity-rotation

```

```

lemmas pf-main-connect = pf-main(1,3,5,7,8-10)[unfolded sr-spectral-radius]
sr-pos[unfolded sr-spectral-radius]
end

```

```

end

```

5.2 Handling Non-Irreducible Matrices as Well

```

theory Perron-Frobenius-General
imports Perron-Frobenius-Irreducible
begin

```

We will need to take sub-matrices and permutations of matrices where the former can best be done via JNF-matrices. So, we first need the Perron-Frobenius theorem in the JNF-world. So, we first define irreducibility of a JNF-matrix.

```

definition graph-of-mat where
graph-of-mat A = (let n = dim-row A; U = {.. $n$ } in
{  $ij. A \ \$ \$ \ ij \neq 0$  }  $\cap U \times U$ )

```

```

definition irreducible-mat where
irreducible-mat A = (let n = dim-row A in
( $\forall i\ j. i < n \longrightarrow j < n \longrightarrow (i,j) \in (graph-of-mat\ A) \hat{+}$ ))

```

```

definition nonneg-irreducible-mat A = (nonneg-mat A  $\wedge$  irreducible-mat A)

```

Next, we have to install transfer rules

```

context
includes lifting-syntax
begin
lemma HMA-irreducible[transfer-rule]: ((HMA-M ::  $- \Rightarrow - \hat{\ } 'n \hat{\ } 'n \Rightarrow -$ )  $====>$ 
(=))
irreducible-mat fixed-mat.irreducible
proof (intro rel-funI, goal-cases)
case (1 a A)
interpret fixed-mat A .
let ?t = Bij-Nat.to-nat :: 'n  $\Rightarrow$  nat
let ?f = Bij-Nat.from-nat :: nat  $\Rightarrow$  'n
from 1[unfolded HMA-M-def]
have a: a = from-hmam A (is - = ?A) by auto
let ?n = CARD('n)

```

```

have dim: dim-row a = ?n unfolding a by simp
have id: {.. $?n$ } = {0.. $?n$ } by auto
have Aij: A $ i $ j = ?A $$ (?t i, ?t j) for i j
  by (metis (no-types, lifting) to-hmam-def to-hma-from-hmam vec-lambda-beta)
have graph: graph-of-mat a =
  {(? $t$  i, ? $t$  j) | i j. A $ i $ j  $\neq$  0} (is ?G = -) unfolding graph-of-mat-def dim
  Let-def id range-to-nat[symmetric]
  unfolding a Aij by auto
have irreducible-mat a = ( $\forall$  i j. i  $\in$  range ?t  $\longrightarrow$  j  $\in$  range ?t  $\longrightarrow$  (i,j)  $\in$  ?G+)

  unfolding irreducible-mat-def dim Let-def range-to-nat by auto
also have ... = ( $\forall$  i j. (?t i, ?t j)  $\in$  ?G+) by auto
also note part1 = calculation
have G: ?G = map-prod ?t ?t ' G unfolding graph G-def by auto
have part2: (?t i, ?t j)  $\in$  ?G+  $\longleftrightarrow$  (i,j)  $\in$  G+ for i j
  unfolding G by (rule inj-trancl-image, simp add: inj-on-def)
show ?case unfolding part1 part2 irreducible-def by auto
qed

```

```

lemma HMA-nonneg-irreducible-mat[transfer-rule]: (HMA-M  $\implies$  (=)) non-
neg-irreducible-mat perron-frobenius
  unfolding perron-frobenius-def pf-nonneg-mat-def perron-frobenius-axioms-def
  nonneg-irreducible-mat-def
  by transfer-prover
end

```

The main statements of Perron-Frobenius can now be transferred to JNF-matrices

```

lemma perron-frobenius-irreducible: fixes A :: real Matrix.mat and cA :: complex
Matrix.mat

```

```

  assumes A: A  $\in$  carrier-mat n n and n: n  $\neq$  0 and nonneg: nonneg-mat A
  and irr: irreducible-mat A
  and cA: cA = map-mat of-real A
  and sr: sr = Spectral-Radius.spectral-radius cA
  shows
    eigenvalue A sr
    order sr (char-poly A) = 1
    0 < sr
    eigenvalue cA  $\alpha \implies$  cmod  $\alpha \leq$  sr
    eigenvalue cA  $\alpha \implies$  cmod  $\alpha =$  sr  $\implies$  order  $\alpha$  (char-poly cA) = 1
     $\exists$  k f. k  $\neq$  0  $\wedge$  k  $\leq$  n  $\wedge$  char-poly A = (monom 1 k - [:sr ^ k:]) * f  $\wedge$ 
      ( $\forall$  x. poly (map-poly complex-of-real f) x = 0  $\longrightarrow$  cmod x < sr)
proof (atomize (full), goal-cases)
  case 1
  from nonneg irr have irr: nonneg-irreducible-mat A unfolding nonneg-irreducible-mat-def
  by auto
  note main = perron-frobenius.pf-main-connect[untransferred, cancel-card-constraint,
  OF A irr,
  folded sr cA]

```

from *main(5,6,7)[OF refl refl n]*
have $\exists k f. k \neq 0 \wedge k \leq n \wedge \text{char-poly } A = (\text{monom } 1 k - [:sr \hat{k}]) * f \wedge$
 $(\forall x. \text{poly } (\text{map-poly complex-of-real } f) x = 0 \longrightarrow \text{cmod } x < sr)$ **by** *blast*
with *main(1,3,8)[OF n] main(2)[OF - n] main(4)[OF - - n]* **show** *?case by auto*
qed

We now need permutations on matrices to show that a matrix if a matrix is not irreducible, then it can be turned into a four-block-matrix by a permutation, where the lower left block is 0.

definition *permutation-mat* :: $\text{nat} \Rightarrow (\text{nat} \Rightarrow \text{nat}) \Rightarrow 'a :: \text{semiring-1 mat}$ **where**
permutation-mat $n p = \text{Matrix.mat } n n (\lambda (i,j). (\text{if } i = p j \text{ then } 1 \text{ else } 0))$

no-notation *m-inv* (*inv1* - [81] 80)

lemma *permutation-mat-dim[simp]*: *permutation-mat* $n p \in \text{carrier-mat } n n$
 $\text{dim-row } (\text{permutation-mat } n p) = n$
 $\text{dim-col } (\text{permutation-mat } n p) = n$
unfolding *permutation-mat-def* **by** *auto*

lemma *permutation-mat-row[simp]*: p *permutes* $\{..<n\} \implies i < n \implies$
 $\text{Matrix.row } (\text{permutation-mat } n p) i = \text{unit-vec } n (\text{inv } p i)$
unfolding *permutation-mat-def unit-vec-def* **by** (*intro eq-vecI, auto simp: permutes-inverses*)

lemma *permutation-mat-col[simp]*: p *permutes* $\{..<n\} \implies i < n \implies$
 $\text{Matrix.col } (\text{permutation-mat } n p) i = \text{unit-vec } n (p i)$
unfolding *permutation-mat-def unit-vec-def* **by** (*intro eq-vecI, auto simp: permutes-inverses*)

lemma *permutation-mat-left*: **assumes** $A: A \in \text{carrier-mat } n nc$ **and** $p: p$ *permutes* $\{..<n\}$

shows $\text{permutation-mat } n p * A = \text{Matrix.mat } n nc (\lambda (i,j). A \ \$\$ (\text{inv } p i, j))$

proof –

{
 fix $i j$
 assume $ij: i < n j < nc$
 from $p \ ij(1)$ **have** $i: \text{inv } p i < n$ **by** (*simp add: permutes-def*)
 have $(\text{permutation-mat } n p * A) \ \$\$ (i,j) = \text{scalar-prod } (\text{unit-vec } n (\text{inv } p i))$
 $(\text{col } A j)$
 by (*subst index-mult-mat, insert ij A p, auto*)
 also have $\dots = A \ \$\$ (\text{inv } p i, j)$
 by (*subst scalar-prod-left-unit, insert A ij i, auto*)
 also note *calculation*
}
thus *?thesis using A*
by (*intro eq-matI, auto*)
qed

lemma permutation-mat-right: **assumes** $A: A \in \text{carrier-mat } nr \ n$ **and** $p: p \text{ permutes } \{..<n\}$
shows $A * \text{permutation-mat } n \ p = \text{Matrix.mat } nr \ n \ (\lambda \ (i,j). \ A \ \$\$ \ (i, \ p \ j))$
proof –
{
 fix $i \ j$
 assume $ij: i < nr \ j < n$
 from $p \ ij(2)$ **have** $j: p \ j < n$ **by** (*simp add: permutes-def*)
 have $(A * \text{permutation-mat } n \ p) \ \$\$ \ (i,j) = \text{scalar-prod } (\text{Matrix.row } A \ i) \ (\text{unit-vec } n \ (p \ j))$
 by (*subst index-mult-mat, insert ij A p, auto*)
 also have $\dots = A \ \$\$ \ (i, \ p \ j)$
 by (*subst scalar-prod-right-unit, insert A ij j, auto*)
 also note *calculation*
}
thus *?thesis* **using** A
 by (*intro eq-matI, auto*)
qed

lemma permutes-lt: $p \text{ permutes } \{..<n\} \implies i < n \implies p \ i < n$
by (*meson lessThan-iff permutes-in-image*)

lemma permutes-iff: $p \text{ permutes } \{..<n\} \implies i < n \implies j < n \implies p \ i = p \ j \longleftrightarrow i = j$
by (*metis permutes-inverses(2)*)

lemma permutation-mat-id-1: **assumes** $p: p \text{ permutes } \{..<n\}$
shows $\text{permutation-mat } n \ p * \text{permutation-mat } n \ (\text{inv } p) = 1_m \ n$
by (*subst permutation-mat-left[OF - p, of - n], force, unfold permutation-mat-def, rule eq-matI,*
auto simp: permutes-lt[OF permutes-inv[OF p]] permutes-iff[OF permutes-inv[OF p]])

lemma permutation-mat-id-2: **assumes** $p: p \text{ permutes } \{..<n\}$
shows $\text{permutation-mat } n \ (\text{inv } p) * \text{permutation-mat } n \ p = 1_m \ n$
by (*subst permutation-mat-right[OF - p, of - n], force, unfold permutation-mat-def, rule eq-matI,*
insert p, auto simp: permutes-lt[OF p] permutes-inverses)

lemma permutation-mat-both: **assumes** $A: A \in \text{carrier-mat } n \ n$ **and** $p: p \text{ permutes } \{..<n\}$
shows $\text{permutation-mat } n \ p * \text{Matrix.mat } n \ n \ (\lambda \ (i,j). \ A \ \$\$ \ (p \ i, \ p \ j)) * \text{permutation-mat } n \ (\text{inv } p) = A$
unfolding *permutation-mat-left[OF mat-carrier p]*
by (*subst permutation-mat-right[OF - permutes-inv[OF p], of - n], force, insert A p,*
auto intro!: eq-matI simp: permutes-inverses permutes-lt[OF permutes-inv[OF p]])

lemma permutation-similar-mat: **assumes** $A: A \in \text{carrier-mat } n \ n$ **and** $p: p \text{ permutes } \{..<n\}$
shows $\text{similar-mat } A \ (\text{Matrix.mat } n \ n \ (\lambda \ (i,j). \ A \ \$\$ \ (p \ i, \ p \ j)))$
by $(\text{rule similar-matI}[\text{OF - permutation-mat-id-1}[\text{OF } p] \ \text{permutation-mat-id-2}[\text{OF } p]]$
 $\text{permutation-mat-both}[\text{symmetric, OF } A \ p], \ \text{insert } A, \ \text{auto})$

lemma det-four-block-mat-lower-left-zero: **fixes** $A1 :: 'a :: \text{idom mat}$
assumes $A1: A1 \in \text{carrier-mat } n \ n$
and $A2: A2 \in \text{carrier-mat } n \ m$ **and** $A30: A3 = 0_m \ m \ n$
and $A4: A4 \in \text{carrier-mat } m \ m$
shows $\text{Determinant.det} \ (\text{four-block-mat } A1 \ A2 \ A3 \ A4) = \text{Determinant.det } A1 \ * \ \text{Determinant.det } A4$
proof –
let $?det = \text{Determinant.det}$
let $?t = \text{transpose-mat}$
let $?A = \text{four-block-mat } A1 \ A2 \ A3 \ A4$
let $?k = n + m$
have $A3: A3 \in \text{carrier-mat } m \ n$ **unfolding** $A30$ **by** auto
have $A: ?A \in \text{carrier-mat } ?k \ ?k$
by $(\text{rule four-block-carrier-mat}[\text{OF } A1 \ A4])$
have $?det \ ?A = ?det \ (?t \ ?A)$
by $(\text{rule sym, rule Determinant.det-transpose}[\text{OF } A])$
also have $?t \ ?A = \text{four-block-mat} \ (?t \ A1) \ (?t \ A3) \ (?t \ A2) \ (?t \ A4)$
by $(\text{rule transpose-four-block-mat}[\text{OF } A1 \ A2 \ A3 \ A4])$
also have $?det \ \dots = ?det \ (?t \ A1) \ * \ ?det \ (?t \ A4)$
by $(\text{rule det-four-block-mat-upper-right-zero}[\text{of - } n \ - \ m], \ \text{insert } A1 \ A2 \ A30 \ A4, \ \text{auto})$
also have $?det \ (?t \ A1) = ?det \ A1$
by $(\text{rule Determinant.det-transpose}[\text{OF } A1])$
also have $?det \ (?t \ A4) = ?det \ A4$
by $(\text{rule Determinant.det-transpose}[\text{OF } A4])$
finally show $?thesis \ .$
qed

lemma char-poly-matrix-four-block-mat: **assumes**
 $A1: A1 \in \text{carrier-mat } n \ n$
and $A2: A2 \in \text{carrier-mat } n \ m$
and $A3: A3 \in \text{carrier-mat } m \ n$
and $A4: A4 \in \text{carrier-mat } m \ m$
shows $\text{char-poly-matrix} \ (\text{four-block-mat } A1 \ A2 \ A3 \ A4) =$
 $\text{four-block-mat} \ (\text{char-poly-matrix } A1) \ (\text{map-mat} \ (\lambda \ x. \ [:-x:]) \ A2)$
 $(\text{map-mat} \ (\lambda \ x. \ [:-x:]) \ A3) \ (\text{char-poly-matrix } A4)$
proof –
from $A1 \ A4$
have $\text{dim}[\text{simp}]: \text{dim-row } A1 = n \ \text{dim-col } A1 = n$
 $\text{dim-row } A4 = m \ \text{dim-col } A4 = m$ **by** auto
show $?thesis$
unfolding $\text{char-poly-matrix-def four-block-mat-def Let-def dim}$

by (rule eq-matI, insert A2 A3, auto)
qed

lemma char-poly-four-block-mat-lower-left-zero: **fixes** A :: 'a :: idom mat
assumes A: A = four-block-mat B C (0_m m n) D
and B: B ∈ carrier-mat n n
and C: C ∈ carrier-mat n m
and D: D ∈ carrier-mat m m
shows char-poly A = char-poly B * char-poly D
unfolding A char-poly-def
by (subst char-poly-matrix-four-block-mat[OF B C - D], force,
rule det-four-block-mat-lower-left-zero[of - n - m], insert B C D, auto)

lemma elements-mat-permutes: **assumes** p: p permutes {.._n}

and A: A ∈ carrier-mat n n

and B: B = Matrix.mat n n (λ (i,j). A \$\$ (p i, p j))

shows elements-mat A = elements-mat B

proof –

from A B **have** [simp]: dim-row A = n dim-col A = n dim-row B = n dim-col
B = n **by** auto

{

fix i j

assume ij: i < n j < n

let ?p = inv p

from permutes-lt[OF p] ij **have** *: p i < n p j < n **by** auto

from permutes-lt[OF permutes-inv[OF p]] ij **have** **: ?p i < n ?p j < n **by**

auto

have ∃ i' j'. B \$\$ (i,j) = A \$\$ (i',j') ∧ i' < n ∧ j' < n

∃ i' j'. A \$\$ (i,j) = B \$\$ (i',j') ∧ i' < n ∧ j' < n

by (rule exI[of - p i], rule exI[of - p j], insert ij *, simp add: B,

rule exI[of - ?p i], rule exI[of - ?p j], insert ** p, simp add: B permutes-inverses)

}

thus ?thesis **unfolding** elements-mat **by** auto

qed

lemma elements-mat-four-block-mat-supseteq:

assumes A1: A1 ∈ carrier-mat n n

and A2: A2 ∈ carrier-mat n m

and A3: A3 ∈ carrier-mat m n

and A4: A4 ∈ carrier-mat m m

shows elements-mat (four-block-mat A1 A2 A3 A4) ⊇

(elements-mat A1 ∪ elements-mat A2 ∪ elements-mat A3 ∪ elements-mat A4)

proof

let ?A = four-block-mat A1 A2 A3 A4

have A: ?A ∈ carrier-mat (n + m) (n + m) **using** A1 A2 A3 A4 **by** simp

from A1 A4

have dim[simp]: dim-row A1 = n dim-col A1 = n

dim-row A4 = m dim-col A4 = m **by** auto

fix x

assume $x: x \in \text{elements-mat } A1 \cup \text{elements-mat } A2 \cup \text{elements-mat } A3 \cup \text{elements-mat } A4$
{
 assume $x \in \text{elements-mat } A1$
 from $\text{this}[\text{unfolded elements-mat}] A1$ **obtain** $i j$ **where** $x: x = A1 \ \$\$ (i,j)$ **and**

 $ij: i < n \ j < n$ **by auto**
 have $x = ?A \ \$\$ (i,j)$ **using** ij **unfolding** x *four-block-mat-def Let-def* **by simp**
 from $\text{elements-matI}[OF A - - \text{this}] ij$ **have** $x \in \text{elements-mat } ?A$ **by auto**
}
moreover
{
 assume $x \in \text{elements-mat } A2$
 from $\text{this}[\text{unfolded elements-mat}] A2$ **obtain** $i j$ **where** $x: x = A2 \ \$\$ (i,j)$ **and**

 $ij: i < n \ j < m$ **by auto**
 have $x = ?A \ \$\$ (i,j + n)$ **using** ij **unfolding** x *four-block-mat-def Let-def* **by simp**
 from $\text{elements-matI}[OF A - - \text{this}] ij$ **have** $x \in \text{elements-mat } ?A$ **by auto**
}
moreover
{
 assume $x \in \text{elements-mat } A3$
 from $\text{this}[\text{unfolded elements-mat}] A3$ **obtain** $i j$ **where** $x: x = A3 \ \$\$ (i,j)$ **and**

 $ij: i < m \ j < n$ **by auto**
 have $x = ?A \ \$\$ (i+n,j)$ **using** ij **unfolding** x *four-block-mat-def Let-def* **by simp**
 from $\text{elements-matI}[OF A - - \text{this}] ij$ **have** $x \in \text{elements-mat } ?A$ **by auto**
}
moreover
{
 assume $x \in \text{elements-mat } A4$
 from $\text{this}[\text{unfolded elements-mat}] A4$ **obtain** $i j$ **where** $x: x = A4 \ \$\$ (i,j)$ **and**

 $ij: i < m \ j < m$ **by auto**
 have $x = ?A \ \$\$ (i+n,j + n)$ **using** ij **unfolding** x *four-block-mat-def Let-def* **by simp**
 from $\text{elements-matI}[OF A - - \text{this}] ij$ **have** $x \in \text{elements-mat } ?A$ **by auto**
}
ultimately show $x \in \text{elements-mat } ?A$ **using** x **by blast**
qed

lemma non-irreducible-mat-split:
fixes $A :: 'a :: \text{idom mat}$
assumes $A: A \in \text{carrier-mat } n \ n$
and not: $\neg \text{irreducible-mat } A$
and $n: n > 1$

shows $\exists n1\ n2\ B\ B1\ B2\ B4. \text{similar-mat } A\ B \wedge \text{elements-mat } A = \text{elements-mat } B \wedge$
 $B = \text{four-block-mat } B1\ B2\ (0_m\ n2\ n1)\ B4 \wedge$
 $B1 \in \text{carrier-mat } n1\ n1 \wedge B2 \in \text{carrier-mat } n1\ n2 \wedge B4 \in \text{carrier-mat } n2\ n2 \wedge$
 $0 < n1 \wedge n1 < n \wedge 0 < n2 \wedge n2 < n \wedge n1 + n2 = n$
proof –
from A **have** $[simp]: \text{dim-row } A = n$ **by** *auto*
let $?G = \text{graph-of-mat } A$
let $?reachp = \lambda\ i\ j. (i,j) \in ?G^{\wedge+}$
let $?reach = \lambda\ i\ j. (i,j) \in ?G^{\wedge*}$
have $\exists\ i\ j. i < n \wedge j < n \wedge \neg ?reach\ i\ j$
proof (*rule ccontr*)
assume $\neg ?thesis$
hence $reach: \bigwedge\ i\ j. i < n \implies j < n \implies ?reach\ i\ j$ **by** *auto*
from $not[\text{unfolded irreducible-mat-def Let-def}]$
obtain $i\ j$ **where** $i: i < n$ **and** $j: j < n$ **and** $nreach: \neg ?reachp\ i\ j$ **by** *auto*
from $reach[OF\ i\ j]\ nreach$ **have** $ij: i = j$ **by** (*simp add: rtrancl-eq-or-trancl*)
from $n\ j$ **obtain** k **where** $k: k < n$ **and** $diff: j \neq k$ **by** *auto*
from $reach[OF\ j\ k]\ diff\ reach[OF\ k\ j]$
have $?reachp\ j\ j$ **by** (*simp add: rtrancl-eq-or-trancl*)
with $nreach\ ij$ **show** *False* **by** *auto*
qed
then obtain $i\ j$ **where** $i: i < n$ **and** $j: j < n$ **and** $nreach: \neg ?reach\ i\ j$ **by** *auto*
define I **where** $I = \{k. k < n \wedge ?reach\ i\ k\}$
have $iI: i \in I$ **unfolding** $I\text{-def}$ **using** $nreach\ i$ **by** *auto*
have $jI: j \notin I$ **unfolding** $I\text{-def}$ **using** $nreach\ j$ **by** *auto*
define f **where** $f = (\lambda\ i. \text{if } i \in I \text{ then } 1 \text{ else } 0 :: \text{nat})$
let $?xs = [0 ..< n]$
from $mset\text{-eq-permutation}[OF\ mset\text{-sort},\ of\ ?xs\ f]$ **obtain** p **where** $p: p$ *permutes*
 $\{..< n\}$
and $perm: \text{permute-list } p\ ?xs = \text{sort-key } f\ ?xs$ **by** *auto*
from p **have** $lt[simp]: i < n \implies p\ i < n$ **for** i **by** (*rule permutes-lt*)
let $?p = \text{inv } p$
have $ip: ?p$ *permutes* $\{..< n\}$ **using** $\text{permutes-inv}[OF\ p]$.
from ip **have** $ilt[simp]: i < n \implies ?p\ i < n$ **for** i **by** (*rule permutes-lt*)
let $?B = \text{Matrix.mat } n\ n\ (\lambda\ (i,j). A\ \$\$ (p\ i,\ p\ j))$
define B **where** $B = ?B$
from $\text{permutation-similar-mat}[OF\ A\ p]$ **have** $sim: \text{similar-mat } A\ B$ **unfolding**
 $B\text{-def}$.
let $?ys = \text{permute-list } p\ ?xs$
define ys **where** $ys = ?ys$
have $len\text{-ys}: \text{length } ys = n$ **unfolding** $ys\text{-def}$ **by** *simp*
let $?k = \text{length } (\text{filter } (\lambda\ i. f\ i = 0)\ ys)$
define k **where** $k = ?k$
have $kn: k \leq n$ **unfolding** $k\text{-def}$ **using** $len\text{-ys}$
using $\text{length-filter-le}[of\ \text{- } ys]$ **by** *auto*
have $ys\text{-p}: i < n \implies ys\ !\ i = p\ i$ **for** i **unfolding** $ys\text{-def permute-list-def}$ **by**
 $simp$

```

have ys: ys = map (λ i. ys ! i) [0 ..< n] unfolding len-ys[symmetric]
  by (simp add: map-nth)
also have ... = map p [0 ..< n]
  by (rule map-cong, insert ys-p, auto)
also have [0 ..< n] = [0 ..< k] @ [k ..< n] using kn
  using le-Suc-ex upt-add-eq-append by blast
finally have ys: ys = map p [0 ..< k] @ map p [k ..< n] by simp
{
  fix i
  assume i: i < n
  let ?g = (λ i. f i = 0)
  let ?f = filter ?g
  from i have pi: p i < n using p by simp
  have k = length (?f ys) by fact
  also have ?f ys = ?f (map p [0 ..< k]) @ ?f (map p [k ..< n]) unfolding ys
by simp
  also note k = calculation
  finally have True by blast
  from perm[symmetric, folded ys-def]
  have sorted (map f ys) using sorted-sort-key by metis
  from this[unfolded ys map-append sorted-append set-map]
  have sorted:  $\bigwedge x y. x < k \implies y \in \{k..<n\} \implies f (p x) \leq f (p y)$  by auto
  have 0:  $\forall i < k. f (p i) = 0$ 
  proof (rule ccontr)
    assume  $\neg$  ?thesis
    then obtain i where i: i < k and zero: f (p i)  $\neq$  0 by auto
    hence f (p i) = 1 unfolding f-def by (auto split: if-splits)
    from sorted[OF i, unfolded this] have 1:  $j \in \{k..<n\} \implies f (p j) \geq 1$  for j
by auto
    have le:  $j \in \{k ..< n\} \implies f (p j) = 1$  for j using 1[of j] unfolding f-def
      by (auto split: if-splits)
    also have ?f (map p [k ..< n]) = [] using le by auto
    from k[unfolded this] have length (?f (map p [0..<k])) = k by simp
    from length-filter-less[of p i map p [0 ..< k] ?g, unfolded this] i zero
    show False by auto
  qed
  hence ?f (map p [0..<k]) = map p [0..<k] by auto
  from arg-cong[OF k[unfolded this, simplified], of set]
  have 1:  $\bigwedge i. i \in \{k ..< n\} \implies f (p i) \neq 0$  by auto
  have 1:  $i < n \implies \neg i < k \implies f (p i) \neq 0$  for i using 1[of i] by auto
  have 0:  $i < n \implies (f (p i) = 0) = (i < k)$  for i using 1[of i] 0[rule-format,
of i] by blast
  have main: (f i = 0) = (?p i < k) using 0[of ?p i] i p
    by (auto simp: permutes-inverses)
  have i ∈ I  $\iff$  f i  $\neq$  0 unfolding f-def by simp
  also have (f i = 0)  $\iff$  ?p i < k using main by auto
  finally have i ∈ I  $\iff$  ?p i  $\geq$  k by auto
} note main = this
from main[OF j] jI

```

```

have k0: k ≠ 0 by auto
from iI main[OF i] have ?p i ≥ k by auto
with ilt[OF i] have kn: k < n by auto
{
  fix i j
  assume i: i < n and ik: k ≤ i and jk: j < k
  with kn have j: j < n by auto
  have jI: p j ∉ I
    by (subst main, insert jk j p, auto simp: permutes-inverses)
  have iI: p i ∈ I
    by (subst main, insert i ik p, auto simp: permutes-inverses)
  from i j have B $$$ (i,j) = A $$$ (p i, p j) unfolding B-def by auto
  also have ... = 0
  proof (rule ccontr)
    assume A $$$ (p i, p j) ≠ 0
    hence (p i, p j) ∈ ?G unfolding graph-of-mat-def Let-def using i j p by
auto
    with iI j have p j ∈ I unfolding I-def by auto
    with jI show False by simp
  qed
  finally have B $$$ (i,j) = 0 .
} note zero = this
have dimB[simp]: dim-row B = n dim-col B = n unfolding B-def by auto
have dim: dim-row B = k + (n - k) dim-col B = k + (n - k) using kn by
auto
obtain B1 B2 B3 B4 where spl: split-block B k k = (B1,B2,B3,B4) (is ?tmp
= -) by (cases ?tmp, auto)
from split-block[OF this dim] have
  Bs: B1 ∈ carrier-mat k k B2 ∈ carrier-mat k (n - k)
  B3 ∈ carrier-mat (n - k) k B4 ∈ carrier-mat (n - k) (n - k)
  and B: B = four-block-mat B1 B2 B3 B4 by auto
have B3: B3 = 0m (n - k) k unfolding arg-cong[OF spl[symmetric], of λ
(-, -, B, -). B, unfolded split]
  unfolding split-block-def Let-def split
  by (rule eq-matI, auto simp: kn zero)
from elements-mat-permutes[OF p A B-def]
have elem: elements-mat A = elements-mat B .
show ?thesis
  by (intro exI conjI, rule sim, rule elem, rule B[unfolded B3], insert Bs k0 kn,
auto)
qed

```

lemma non-irreducible-nonneg-mat-split:

```

fixes A :: 'a :: linordered-idom mat
assumes A: A ∈ carrier-mat n n
and nonneg: nonneg-mat A
and not: ¬ irreducible-mat A
and n: n > 1
shows ∃ n1 n2 A1 A2. char-poly A = char-poly A1 * char-poly A2

```

```

    ∧ nonneg-mat A1 ∧ nonneg-mat A2
    ∧ A1 ∈ carrier-mat n1 n1 ∧ A2 ∈ carrier-mat n2 n2
    ∧ 0 < n1 ∧ n1 < n ∧ 0 < n2 ∧ n2 < n ∧ n1 + n2 = n
proof –
  from non-irreducible-mat-split[OF A not n]
  obtain n1 n2 B B1 B2 B4
    where sim: similar-mat A B and elem: elements-mat A = elements-mat B
    and B: B = four-block-mat B1 B2 (0m n2 n1) B4
    and Bs: B1 ∈ carrier-mat n1 n1 B2 ∈ carrier-mat n1 n2 B4 ∈ carrier-mat
n2 n2
    and n: 0 < n1 n1 < n 0 < n2 n2 < n n1 + n2 = n by auto
  from char-poly-similar[OF sim]
  have AB: char-poly A = char-poly B .
  from nonneg have nonneg: nonneg-mat B unfolding nonneg-mat-def elem by
auto
  have cB: char-poly B = char-poly B1 * char-poly B4
    by (rule char-poly-four-block-mat-lower-left-zero[OF B Bs])
  from nonneg have B1-B4: nonneg-mat B1 nonneg-mat B4 unfolding B non-
neg-mat-def
    using elements-mat-four-block-mat-supseteq[OF Bs(1–2) - Bs(3), of 0m n2
n1] by auto
  show ?thesis
    by (intro exI conjI, rule AB[unfolded cB], rule B1-B4, rule B1-B4,
rule Bs, rule Bs, insert n, auto)
qed

```

The main generalized theorem. The characteristic polynomial of a non-negative real matrix can be represented as a product of roots of unitys (scaled by the the spectral radius sr) and a polynomial where all roots are smaller than the spectral radius.

```

theorem perron-frobenius-nonneg: fixes A :: real Matrix.mat
assumes A: A ∈ carrier-mat n n and pos: nonneg-mat A and n: n ≠ 0
shows ∃ sr ks f.
  sr ≥ 0 ∧
  0 ∉ set ks ∧ ks ≠ [] ∧
  char-poly A = prod-list (map (λ k. monom 1 k - [:sr ^ k:]) ks) * f ∧
  (∀ x. poly (map-poly complex-of-real f) x = 0 → cmod x < sr)
proof –
  define p where p = (λ sr k. monom 1 k - [(sr :: real) ^ k:])
  let ?small = λ f sr. (∀ x. poly (map-poly complex-of-real f) x = 0 → cmod x
< sr)
  let ?wit = λ A sr ks f. sr ≥ 0 ∧ 0 ∉ set ks ∧ ks ≠ [] ∧
char-poly A = prod-list (map (p sr) ks) * f ∧ ?small f sr
  let ?c = complex-of-real
  interpret c: field-hom ?c ..
  interpret p: map-poly-inj-idom-divide-hom ?c ..
  have map-p: map-poly ?c (p sr k) = (monom 1 k - [:?c sr ^ k:]) for sr k
    unfolding p-def by (simp add: hom-distrib)
  {

```

```

fix k x sr
assume 0: poly (map-poly ?c (p sr k)) x = 0 and k: k ≠ 0 and sr: sr ≥ 0
note 0 also note map-p
finally have x̂k = (?c sr)̂k by (simp add: poly-monom)
from arg-cong[OF this, of λ c. root k (cmod c), unfolded norm-power] k
have cmod x = cmod (?c sr) using real-root-pos2 by auto
also have ... = sr using sr by auto
finally have cmod x = sr .
} note p-conv = this
have ∃ sr ks f. ?wit A sr ks f using A pos n
proof (induct n arbitrary: A rule: less-induct)
  case (less n A)
  note pos = less(3)
  note A = less(2)
  note IH = less(1)
  note n = less(4)
  from n
  consider (1) n = 1
    | (irr) irreducible-mat A
    | (red) ¬ irreducible-mat A n > 1
  by force
  thus ∃ sr ks f. ?wit A sr ks f
  proof cases
    case irr
    from perron-frobenius-irreducible(3,6)[OF A n pos irr refl refl]
    obtain sr k f where
      *: sr > 0 k ≠ 0 char-poly A = p sr k * f ?small f sr unfolding p-def
    by auto
    hence ?wit A sr [k] f by auto
    thus ?thesis by blast
  next
  case red
  from non-irreducible-nonneg-mat-split[OF A pos red] obtain n1 n2 A1 A2
  where char: char-poly A = char-poly A1 * char-poly A2
    and pos: nonneg-mat A1 nonneg-mat A2
    and A: A1 ∈ carrier-mat n1 n1 A2 ∈ carrier-mat n2 n2
    and n: n1 < n n2 < n
    and n0: n1 ≠ 0 n2 ≠ 0 by auto
  from IH[OF n(1) A(1) pos(1) n0(1)] obtain sr1 ks1 f1 where 1: ?wit A1
sr1 ks1 f1 by blast
  from IH[OF n(2) A(2) pos(2) n0(2)] obtain sr2 ks2 f2 where 2: ?wit A2
sr2 ks2 f2 by blast
  have ∃ A1 A2 sr1 ks1 f1 sr2 ks2 f2. ?wit A1 sr1 ks1 f1 ∧ ?wit A2 sr2 ks2 f2
  ∧
    sr1 ≥ sr2 ∧ char-poly A = char-poly A1 * char-poly A2
  proof (cases sr1 ≥ sr2)
    case True
    show ?thesis unfolding char
    by (intro exI, rule conjI[OF 1 conjI[OF 2]], insert True, auto)
  end

```

```

next
  case False
  show ?thesis unfolding char
    by (intro exI, rule conjI[OF 2 conjI[OF 1]], insert False, auto)
qed
then obtain A1 A2 sr1 ks1 f1 sr2 ks2 f2 where
  1: ?wit A1 sr1 ks1 f1 and 2: ?wit A2 sr2 ks2 f2 and
  sr: sr1 ≥ sr2 and char: char-poly A = char-poly A1 * char-poly A2 by
blast
show ?thesis
proof (cases sr1 = sr2)
  case True
  have ?wit A sr2 (ks1 @ ks2) (f1 * f2) unfolding char
    by (insert 1 2 True, auto simp: True p.hom-mult)
  thus ?thesis by blast
next
  case False
  with sr have sr1: sr1 > sr2 by auto
  have lt: poly (map-poly ?c (p sr2 k)) x = 0 ⇒ k ∈ set ks2 ⇒ cmod x <
sr1 for k x
    using sr1 p-conv[of sr2 k x] 2 by auto
  have ?wit A sr1 ks1 (f1 * f2 * prod-list (map (p sr2) ks2)) unfolding char
    by (insert 1 2 sr1 lt, auto simp: p.hom-mult p.hom-prod-list
      poly-prod-list prod-list-zero-iff)
  thus ?thesis by blast
qed
next
  case 1
  define a where a = A $$ (0,0)
  have A: A = Matrix.mat 1 1 (λ x. a)
    by (rule eq-matI, unfold a-def, insert A 1(1), auto)
  have char: char-poly A = [: - a, 1 :] unfolding A
    by (auto simp: Determinant.det-def char-poly-def char-poly-matrix-def)
  from pos A have a: a ≥ 0 unfolding nonneg-mat-def elements-mat by auto
  have ?wit A a [1] 1 unfolding char using a by (auto simp: p-def monom-Suc)
  thus ?thesis by blast
qed
qed
then obtain sr ks f where wit: ?wit A sr ks f by blast
thus ?thesis using wit unfolding p-def by auto
qed

```

And back to HMA world via transfer.

```

theorem perron-frobenius-non-neg: fixes A :: real ^ 'n ^ 'n
  assumes pos: non-neg-mat A
  shows  $\exists sr\ ks\ f.$ 
     $sr \geq 0 \wedge$ 
     $0 \notin set\ ks \wedge ks \neq [] \wedge$ 
     $charpoly\ A = prod-list\ (map\ (\lambda k. monom\ 1\ k - [:sr \wedge k:] ks) * f \wedge$ 

```

```

  (∀ x. poly (map-poly complex-of-real f) x = 0 → cmod x < sr)
using pos
proof (transfer, goal-cases)
  case (1 A)
  from perron-frobenius-nonneg[OF 1]
  show ?case by auto
qed

```

We now specialize the theorem for complexity analysis where we are mainly interested in the case where the spectral radius is at most 1. Note that this can be checked by testing that there are no real roots of the characteristic polynomial which exceed 1.

Moreover, here the existential quantifier over the factorization is replaced by *decompose-prod-root-unity*, an algorithm which computes this factorization in an efficient way.

lemma *perron-frobenius-for-complexity*: **fixes** $A :: \text{real}^{\wedge n} \wedge \text{real}^{\wedge n}$ **and** $f :: \text{real poly}$

```

defines cA ≡ map-matrix complex-of-real A
defines cf ≡ map-poly complex-of-real f
assumes pos: non-neg-mat A
and sr: ∧ x. poly (charpoly A) x = 0 ⇒ x ≤ 1
and decomp: decompose-prod-root-unity (charpoly A) = (ks, f)
shows 0 ∉ set ks
  charpoly A = prod-root-unity ks * f
  charpoly cA = prod-root-unity ks * cf
  ∧ x. poly (charpoly cA) x = 0 ⇒ cmod x ≤ 1
  ∧ x. poly cf x = 0 ⇒ cmod x < 1
  ∧ x. cmod x = 1 ⇒ order x (charpoly cA) = length [k←ks . xk = 1]
  ∧ x. cmod x = 1 ⇒ poly (charpoly cA) x = 0 ⇒ ∃ k ∈ set ks. xk = 1
unfolding cf-def cA-def
proof (atomize(full), goal-cases)
  case 1
  let ?c = complex-of-real
  let ?cp = map-poly ?c
  let ?A = map-matrix ?c A
  let ?wit = λ ks f. 0 ∉ set ks ∧
    charpoly A = prod-root-unity ks * f ∧
    charpoly ?A = prod-root-unity ks * map-poly of-real f ∧
    (∀ x. poly (charpoly ?A) x = 0 → cmod x ≤ 1) ∧
    (∀ x. poly (?cp f) x = 0 → cmod x < 1)
  interpret field-hom ?c ..
  interpret p: map-poly-inj-idom-divide-hom ?c ..
  {
    from perron-frobenius-non-neg[OF pos] obtain sr ks f
    where *: sr ≥ 0 0 ∉ set ks ks ≠ []
    and cp: charpoly A = prod-list (map (λ k. monom 1 k - [:srk:]) ks) * f
    and small: ∧ x. poly (?cp f) x = 0 ⇒ cmod x < sr by blast

    from arg-cong[OF cp, of map-poly ?c]
  }

```

```

have cpc: charpoly ?A = prod-list (map (λ k. monom 1 k - [?:c sr ^ k:]) ks) *
map-poly ?c f
by (simp add: charpoly-of-real hom-distrib p.prod-list-map-hom[symmetric]
o-def)
have sr-le-1: sr ≤ 1
by (rule sr, unfold cp, insert *, cases ks, auto simp: poly-monom)
{
fix x
note [simp] = prod-list-zero-iff o-def poly-monom
assume poly (charpoly ?A) x = 0
from this[unfolded cpc poly-mult poly-prod-list] small[of x]
consider (lt) cmod x < sr | (mem) k where k ∈ set ks x ^ k = (?c sr) ^ k
by force
hence cmod x ≤ sr
proof (cases)
case (mem k)
with * have k: k ≠ 0 by metis
with arg-cong[OF mem(2), of λ x. root k (cmod x), unfolded norm-power]
real-root-pos2[of k] *(1)
have cmod x = sr by auto
thus ?thesis by auto
qed simp
} note root = this
have ∃ ks f. ?wit ks f
proof (cases sr = 1)
case False
with sr-le-1 have *: cmod x ≤ sr ⇒ cmod x < 1 cmod x ≤ sr ⇒ cmod x
≤ 1 for x by auto
show ?thesis
by (rule exI[of - Nil], rule exI[of - charpoly A], insert * root,
auto simp: prod-root-unity-def charpoly-of-real)
next
case sr: True
from * cp cpc small root
show ?thesis unfolding sr root-unity-def prod-root-unity-def by (auto simp:
pCons-one)
qed
}
then obtain Ks F where wit: ?wit Ks F by auto
have cA0: charpoly ?A ≠ 0 using degree-monic-charpoly[of ?A] by auto
from wit have id: charpoly ?A = prod-root-unity Ks * ?cp F by auto
from of-real-hom.hom-decompose-prod-root-unity[of charpoly A, unfolded decomp]
have decompc: decompose-prod-root-unity (charpoly ?A) = (ks, ?cp f)
by (auto simp: charpoly-of-real)
from wit have small: cmod x = 1 ⇒ poly (?cp F) x ≠ 0 for x by auto
from decompose-prod-root-unity[OF id decompc this cA0]
have id: charpoly ?A = prod-root-unity ks * ?cp F F = f set Ks by auto
have ?cp (charpoly A) = ?cp (prod-root-unity ks * f) unfolding id
unfolding charpoly-of-real[symmetric] id p.hom-mult of-real-hom.hom-prod-root-unity

```



```

..
hence idr: charpoly  $A = \text{prod-root-unity } ks * f$  by auto
have wit:  $?wit\ ks\ f$  and idc: charpoly  $?A = \text{prod-root-unity } ks * ?cp\ f$ 
using wit unfolding id idr by auto
{
  fix x
  assume cmod  $x = 1$ 
  from small[OF this, unfolded id] have poly  $(?cp\ f)\ x \neq 0$  by auto
  from order-OI[OF this] this have ord: order  $x\ (?cp\ f) = 0$  and cf0:  $?cp\ f \neq 0$ 
by auto
  have order  $x\ (\text{charpoly } ?A) = \text{order } x\ (\text{prod-root-unity } ks)$  unfolding idc
    by (subst order-mult, insert cf0 wit ord, auto)
  also have  $\dots = \text{length } [k \leftarrow ks . x \wedge k = 1]$ 
    by (subst order-prod-root-unity, insert wit, auto)
  finally have ord: order  $x\ (\text{charpoly } ?A) = \text{length } [k \leftarrow ks . x \wedge k = 1]$  .
  {
    assume poly  $(\text{charpoly } ?A)\ x = 0$ 
    with cA0 have order  $x\ (\text{charpoly } ?A) \neq 0$  unfolding order-root by auto
    from this[unfolded ord] have  $\exists k \in \text{set } ks . x \wedge k = 1$ 
    by (cases  $[k \leftarrow ks . x \wedge k = 1]$ , force+)
  }
  note this ord
}
with wit show ?case by blast
qed

```

and convert to JNF-world

```

lemmas perron-frobenius-for-complexity-jnf =
  perron-frobenius-for-complexity[unfolded atomize-imp atomize-all,
    untransferred, cancel-card-constraint, rule-format]

```

end

6 Combining Spectral Radius Theory with Perron Frobenius theorem

```

theory Spectral-Radius-Theory

```

```

imports

```

```

  Polynomial-Factorization.Square-Free-Factorization

```

```

  Jordan-Normal-Form.Spectral-Radius

```

```

  Jordan-Normal-Form.Char-Poly

```

```

  Perron-Frobenius

```

```

  HOL-Computational-Algebra.Field-as-Ring

```

```

begin

```

```

abbreviation spectral-radius where spectral-radius  $\equiv \text{Spectral-Radius.spectral-radius}$ 

```

```

hide-const (open) Module.smult

```

Via JNFs it has been proven that the growth of A^k is polynomially bounded, if all complex eigenvalues have a norm at most 1, i.e., the spectral

radius must be at most 1. Moreover, the degree of the polynomial growth can be bounded by the order of those roots which have norm 1, cf. $\llbracket ?A \in \text{carrier-mat } ?n \text{ } ?n; \text{Spectral-Radius-Theory.spectral-radius } ?A \leq 1; \bigwedge \text{ev } k. \llbracket \text{poly } (\text{char-poly } ?A) \text{ ev} = 0; \text{cmod ev} = 1 \rrbracket \implies \text{order ev } (\text{char-poly } ?A) \leq ?d \rrbracket \implies \exists c1 \ c2. \forall k. \text{norm-bound } (?A \hat{=} k) (c1 + c2 * (\text{real } k)^{?d - 1})$.

Perron Frobenius theorem tells us that for a real valued non negative matrix, the largest eigenvalue is a real non-negative one. Hence, we only have to check, that all real eigenvalues are at most one.

We combine both theorems in the following. To be more precise, the set-based complexity results from JNFs with the type-based Perron Frobenius theorem in HMA are connected to obtain a set based complexity criterion for real-valued non-negative matrices, where one only investigated the real valued eigenvalues for checking the eigenvalue-at-most-1 condition. Here, in the precondition of the roots of the polynomial, the type-system ensures that we only have to look at real-valued eigenvalues, and can ignore the complex-valued ones.

The linkage between set-and type-based is performed via HMA-connect.

lemma *perron-frobenius-spectral-radius-complex*: **fixes** $A :: \text{complex mat}$
assumes $A: A \in \text{carrier-mat } n \ n$
and *real-nonneg*: $\text{real-nonneg-mat } A$
and *ev-le-1*: $\bigwedge x. \text{poly } (\text{char-poly } (\text{map-mat } \text{Re } A)) \ x = 0 \implies x \leq 1$
and *ev-order*: $\bigwedge x. \text{norm } x = 1 \implies \text{order } x \ (\text{char-poly } A) \leq d$
shows $\exists c1 \ c2. \forall k. \text{norm-bound } (A \hat{=} k) (c1 + c2 * \text{real } k \wedge (d - 1))$
proof (*cases* $n = 0$)
case *False*
hence $n: n > 0 \ n \neq 0$ **by** *auto*
define *sr* **where** $sr = \text{spectral-radius } A$
note $sr = \text{spectral-radius-mem-max}[OF \ A \ n(1), \text{folded } sr\text{-def}]$
show *thesis*
proof (*rule* *spectral-radius-poly-bound*[*OF* A], *unfold* $sr\text{-def}$ [*symmetric*])
let $?cr = \text{complex-of-real}$

here is the transition from type-based perron-frobenius to set-based
from *perron-frobenius*[*untransferred*, *cancel-card-constraint*, *OF* A *real-nonneg* $n(2)$]
obtain v **where** $v: v \in \text{carrier-vec } n$ **and** $ev: \text{eigenvector } A \ v \ (?cr \ sr)$ **and**
 $rnn: \text{real-nonneg-vec } v$ **unfolding** $sr\text{-def}$ **by** *auto*
define B **where** $B = \text{map-mat } \text{Re } A$
let $?A = \text{map-mat } ?cr \ B$
have $AB: A = ?A$ **unfolding** $B\text{-def}$
by (*rule* *eq-matI*, *insert* *real-nonneg*[*unfolded* *real-nonneg-mat-def* *elements-mat-def*], *auto*)
define w **where** $w = \text{map-vec } \text{Re } v$
let $?v = \text{map-vec } ?cr \ w$
have $vw: v = ?v$ **unfolding** $w\text{-def}$

```

    by (rule eq-vecI, insert rnn[unfolded real-nonneg-vec-def vec-elements-def],
    auto)
  have B: B ∈ carrier-mat n n unfolding B-def using A by auto
  from AB vw ev have ev: eigenvector ?A ?v (?cr sr) by simp
  have eigenvector B w sr
    by (rule of-real-hom.eigenvector-hom-vec[OF B ev])
  hence eigenvalue B sr unfolding eigenvalue-def by blast
  from ev-le-1[folded B-def, OF this[unfolded eigenvalue-root-char-poly[OF B]]]
  show sr ≤ 1 .
next
fix ev
assume cmod ev = 1
thus order ev (char-poly A) ≤ d by (rule ev-order)
qed
next
case True
with A show ?thesis
  by (intro exI[of - 0], auto simp: norm-bound-def)
qed

```

The following lemma is the same as $\llbracket ?A \in \text{carrier-mat } ?n \ ?n; \text{real-nonneg-mat } ?A; \bigwedge x. \text{poly}(\text{char-poly}(\text{map-mat Re } ?A)) x = 0 \implies x \leq 1; \bigwedge x. \text{cmod } x = 1 \implies \text{order } x(\text{char-poly } ?A) \leq ?d \rrbracket \implies \exists c1 \ c2. \forall k. \text{norm-bound} (?A \hat{\ }_m k) (c1 + c2 * (\text{real } k)^{d-1})$, except that now the type *real* is used instead of *complex*.

```

lemma perron-frobenius-spectral-radius: fixes A :: real mat
assumes A: A ∈ carrier-mat n n
and nonneg: nonneg-mat A
and ev-le-1:  $\forall x. \text{poly}(\text{char-poly } A) x = 0 \implies x \leq 1$ 
and ev-order:  $\forall x :: \text{complex}. \text{norm } x = 1 \implies \text{order } x(\text{map-poly of-real}(\text{char-poly } A)) \leq d$ 
shows  $\exists c1 \ c2. \forall k \ a. a \in \text{elements-mat}(A \hat{\ }_m k) \implies \text{abs } a \leq (c1 + c2 * \text{real } k)^{d-1}$ 
proof -
  let ?cr = complex-of-real
  let ?B = map-mat ?cr A
  have B: ?B ∈ carrier-mat n n using A by auto
  have rnn: real-nonneg-mat ?B using nonneg unfolding real-nonneg-mat-def
  nonneg-mat-def
    by (auto simp: elements-mat-def)
  have id: map-mat Re ?B = A
    by (rule eq-matI, auto)
  have  $\exists c1 \ c2. \forall k. \text{norm-bound} (?B \hat{\ }_m k) (c1 + c2 * \text{real } k)^{d-1}$ 
    by (rule perron-frobenius-spectral-radius-complex[OF B rnn], unfold id,
    insert ev-le-1 ev-order, auto simp: of-real-hom.char-poly-hom[OF A])
  then obtain c1 c2 where nb:  $\bigwedge k. \text{norm-bound} (?B \hat{\ }_m k) (c1 + c2 * \text{real } k)^{d-1}$ 
  by auto
  show ?thesis
  proof (rule exI[of - c1], rule exI[of - c2], intro allI impI)

```

fix $k a$
assume $a \in \text{elements-mat } (A \hat{\ }_m k)$
with $\text{pow-carrier-mat}[OF A]$ **obtain** ij **where** $a: a = (A \hat{\ }_m k) \ \$\$ (i,j)$ **and**
 $ij: i < n \ j < n$
unfolding elements-mat **by force**
from $ij \text{ nb}[of k] A$ **have** $\text{norm } ((?B \hat{\ }_m k) \ \$\$ (i,j)) \leq c1 + c2 * \text{real } k \hat{\ }^{(d - 1)}$
unfolding norm-bound-def **by auto**
also have $(?B \hat{\ }_m k) \ \$\$ (i,j) = ?cr a$
unfolding $\text{of-real-hom.mat-hom-pow}[OF A, \text{symmetric}] a$ **using** $ij A$ **by auto**
also have $\text{norm } (?cr a) = \text{abs } a$ **by auto**
finally show $\text{abs } a \leq (c1 + c2 * \text{real } k \hat{\ }^{(d - 1)})$.
qed
qed

We can also convert the set-based lemma $\llbracket ?A \in \text{carrier-mat } ?n \ ?n; \text{nonneg-mat } ?A; \forall x. \text{poly } (\text{char-poly } ?A) \ x = 0 \longrightarrow x \leq 1; \forall x. \text{cmod } x = 1 \longrightarrow \text{order } x \ (\text{map-poly } \text{complex-of-real } (\text{char-poly } ?A)) \leq ?d \rrbracket \implies \exists c1 \ c2. \forall k \ a. a \in \text{elements-mat } (?A \hat{\ }_m k) \longrightarrow |a| \leq c1 + c2 * (\text{real } k)^{?d - 1}$ to a type-based version.

lemma $\text{perron-frobenius-spectral-type-based}$:
assumes $\text{non-neg-mat } (A :: \text{real } \hat{\ }^n \hat{\ }^n)$
and $\forall x. \text{poly } (\text{charpoly } A) \ x = 0 \longrightarrow x \leq 1$
and $\forall x :: \text{complex. norm } x = 1 \longrightarrow \text{order } x \ (\text{map-poly of-real } (\text{charpoly } A)) \leq d$
shows $\exists c1 \ c2. \forall k \ a. a \in \text{elements-mat-h } (\text{matpow } A \ k) \longrightarrow \text{abs } a \leq (c1 + c2 * \text{real } k \hat{\ }^{(d - 1)})$
using $\text{assms } \text{perron-frobenius-spectral-radius}$
by $(\text{transfer}, \text{blast})$

And of course, we can also transfer the type-based lemma back to a set-based setting, only that – without further case-analysis – we get the additional assumption $n \neq 0$.

lemma **assumes** $A \in \text{carrier-mat } n \ n$
and $\text{nonneg-mat } A$
and $\forall x. \text{poly } (\text{char-poly } A) \ x = 0 \longrightarrow x \leq 1$
and $\forall x :: \text{complex. norm } x = 1 \longrightarrow \text{order } x \ (\text{map-poly of-real } (\text{char-poly } A)) \leq d$
and $n \neq 0$
shows $\exists c1 \ c2. \forall k \ a. a \in \text{elements-mat } (A \hat{\ }_m k) \longrightarrow \text{abs } a \leq (c1 + c2 * \text{real } k \hat{\ }^{(d - 1)})$
using $\text{perron-frobenius-spectral-type-based}[\text{untransferred}, \text{cancel-card-constraint}, \text{OF assms}]$.

Note that the precondition eigenvalue-at-most-1 can easily be formulated as a cardinality constraints which can be decided by Sturm’s theorem. And in order to obtain a bound on the order, one can perform a square-free-factorization (via Yun’s factorization algorithm) of the characteristic polynomial into $f_1^1 \cdot \dots \cdot f_d^d$ where each f_i has precisely the roots of order i .

context

fixes $A :: \text{real mat}$ **and** $c :: \text{real}$ **and** fis **and** $n :: \text{nat}$
assumes $A: A \in \text{carrier-mat } n \ n$
and $\text{nonneg}: \text{nonneg-mat } A$
and $\text{yun}: \text{yun-factorization gcd } (\text{char-poly } A) = (c, fis)$
and $\text{ev-le-1}: \text{card } \{x. \text{poly } (\text{char-poly } A) \ x = 0 \wedge x > 1\} = 0$

begin

Note that *yun-factorization* has an offset by 1, so the pair $(f_i, i) \in \text{set } fis$ encodes $f_i^{\text{Suc } i}$.

lemma *perron-frobenius-spectral-radius-yun*:

assumes $\text{bnd}: \bigwedge f_i \ i. (f_i, i) \in \text{set } fis$
 $\implies (\exists x :: \text{complex}. \text{poly } (\text{map-poly of-real } f_i) \ x = 0 \wedge \text{norm } x = 1)$
 $\implies \text{Suc } i \leq d$

shows $\exists c1 \ c2. \forall k \ a. a \in \text{elements-mat } (A \widehat{\ }_m \ k) \longrightarrow \text{abs } a \leq (c1 + c2 * \text{real } k \wedge (d - 1))$

proof (*rule perron-frobenius-spectral-radius[OF A nonneg]; intro allI impI*)

let $?cr = \text{complex-of-real}$

let $?cp = \text{map-poly } ?cr \ (\text{char-poly } A)$

fix $x :: \text{complex}$

assume $x: \text{norm } x = 1$

have $A0: \text{char-poly } A \neq 0$ **using** *degree-monic-char-poly[OF A]* **by** *auto*

interpret *field-hom-0' ?cr* **by** (*standard, auto*)

from $A0$ **have** $cp0: ?cp \neq 0$ **by** *auto*

obtain ox **where** $ox: \text{order } x \ ?cp = ox$ **by** *blast*

note $\text{sff} = \text{square-free-factorization-order-root}[OF \ \text{yun-factorization}(1)[OF \ \text{yun-factorization-hom}[of \ \text{char-poly } A, \ \text{unfolded yun map-prod-def split}]] \ cp0, \ of \ x \ ox, \ \text{unfolded } ox]$

show $\text{order } x \ ?cp \leq d$ **unfolding** ox

proof (*cases ox*)

case ($\text{Suc } oo$)

with sff **obtain** fi **where** $\text{mem}: (fi, oo) \in \text{set } fis$ **and** $rt: \text{poly } (\text{map-poly } ?cr \ fi) \ x = 0$ **by** *auto*

from $\text{bnd}[OF \ \text{mem } \text{exI}[of \ - \ x], \ OF \ \text{conjI}[OF \ rt \ x]]$

show $ox \leq d$ **unfolding** Suc .

qed *auto*

next

let $?L = \{x. \text{poly } (\text{char-poly } A) \ x = 0 \wedge x > 1\}$

fix $x :: \text{real}$

assume $rt: \text{poly } (\text{char-poly } A) \ x = 0$

have $\text{finite } ?L$

by (*rule finite-subset[OF - poly-roots-finite[of char-poly A]], insert degree-monic-char-poly[OF A], auto*)

with *ev-le-1* **have** $?L = \{\}$ **by** *simp*

with rt **show** $x \leq 1$ **by** *auto*

qed

Note that the only remaining problem in applying $(\bigwedge f_i \ i. \llbracket (f_i, i) \in \text{set } fis; \exists x. \text{poly } (\text{map-poly complex-of-real } f_i) \ x = 0 \wedge \text{cmod } x = 1 \rrbracket \implies \text{Suc } i \leq ?d) \implies \exists c1 \ c2. \forall k \ a. a \in \text{elements-mat } (A \widehat{\ }_m \ k) \longrightarrow |a| \leq c1 + c2 *$

$(real\ k)^{?d-1}$ is to check the condition $\exists x. poly\ (map\ poly\ complex\ of\ real\ f_i)\ x = 0 \wedge cmod\ x = 1$. Here, there are at least three possibilities. First, one can just ignore this precondition and weaken the statement. Second, one can apply Sturm's theorem to determine whether all roots are real. This can be done by comparing the number of distinct real roots with the degree of f_i , since f_i is square-free. If all roots are real, then one can decide the criterion by checking the only two possible real roots with norm equal to 1, namely 1 and -1. If on the other hand there are complex roots, then we loose precision at this point. Third, one uses a factorization algorithm (e.g., via complex algebraic numbers) to precisely determine the complex roots and decide the condition.

The second approach is illustrated in the following theorem. Note that all preconditions – including the ones from the context – can easily be checked with the help of Sturm's method. This method is used as a fast approximative technique in CeTA [3]. Only if the desired degree cannot be ensured by this method, the more costly complex algebraic number based factorization is applied.

lemma *perron-frobenius-spectral-radius-yun-real-roots:*

assumes *bnd*: $\bigwedge f_i\ i. (f_i, i) \in set\ fis$

$\implies card\ \{x. poly\ f_i\ x = 0\} \neq degree\ f_i \vee poly\ f_i\ 1 = 0 \vee poly\ f_i\ (-1) = 0$

$\implies Suc\ i \leq d$

shows $\exists c1\ c2. \forall k\ a. a \in elements\ mat\ (A\ \widehat{m}\ k) \implies abs\ a \leq (c1 + c2 * real\ k\ \widehat{(d-1)})$

proof (*rule perron-frobenius-spectral-radius-yun*)

fix $f_i\ i$

let $?cr = complex\ of\ real$

let $?cp = map\ poly\ ?cr$

assume $fi: (f_i, i) \in set\ fis$

and $\exists x. poly\ (map\ poly\ ?cr\ fi)\ x = 0 \wedge norm\ x = 1$

then obtain x **where** $rt: poly\ (?cp\ fi)\ x = 0$ **and** $x: norm\ x = 1$ **by** *auto*

show $Suc\ i \leq d$

proof (*rule bnd[OF fi]*)

consider $(c)\ x \notin \mathbb{R} \mid (1)\ x = 1 \mid (m1)\ x = -1 \mid (r)\ x \in \mathbb{R}\ x \notin \{1, -1\}$

by (*cases* $x \in \mathbb{R}$; *auto*)

thus $card\ \{x. poly\ f_i\ x = 0\} \neq degree\ f_i \vee poly\ f_i\ 1 = 0 \vee poly\ f_i\ (-1) = 0$

proof (*cases*)

case *1*

from rt **have** $poly\ f_i\ 1 = 0$

unfolding *1* **by** *simp*

thus $?thesis$ **by** *simp*

next

case *m1*

have $id: -1 = ?cr\ (-1)$ **by** *simp*

from rt **have** $poly\ f_i\ (-1) = 0$

unfolding *m1 id of-real-hom.hom-zero* **where** $'a = complex, symmetric$

of-real-hom.poly-map-poly **by** *simp*

thus $?thesis$ **by** *simp*

```

next
  case r
  then obtain y where xy: x = of-real y unfolding Reals-def by auto
  from r(2)[unfolded xy] have y: y  $\notin$  {1, -1} by auto
  from x[unfolded xy] have abs y = 1 by auto
  with y have False by auto
  thus ?thesis ..
next
  case c
  from yun-factorization(2)[OF yun] fi have monic fi by auto
  hence fi: ?cp fi  $\neq$  0 by auto
  hence fin: finite {x. poly (?cp fi) x = 0} by (rule poly-roots-finite)
  have ?cr ' {x. poly (?cp fi) (?cr x) = 0}  $\subset$  {x. poly (?cp fi) x = 0} (is ?l  $\subset$ 
?r)
  proof (rule, force)
    have x  $\in$  ?r using rt by auto
    moreover have x  $\notin$  ?l using c unfolding Reals-def by auto
    ultimately show ?l  $\neq$  ?r by blast
  qed
  from psubset-card-mono[OF fin this] have card ?l < card ?r .
  also have ...  $\leq$  degree (?cp fi) by (rule poly-roots-degree[OF fi])
  also have ... = degree fi by simp
  also have ?l = ?cr ' {x. poly fi x = 0} by auto
  also have card ... = card {x. poly fi x = 0}
    by (rule card-image, auto simp: inj-on-def)
  finally have card {x. poly fi x = 0}  $\neq$  degree fi by simp
  thus ?thesis by auto
qed
qed
qed
end
end

```

7 The Jordan Blocks of the Spectral Radius are Largest

Consider a non-negative real matrix, and consider any Jordan-block of any eigenvalues whose norm is the spectral radius. We prove that there is a Jordan block of the spectral radius which has the same size or is larger.

theory *Spectral-Radius-Largest-Jordan-Block*

imports

Jordan-Normal-Form.Jordan-Normal-Form-Uniqueness

Perron-Frobenius-General

HOL-Real-Asymp.Real-Asymp

begin

lemma *poly-asymp-equiv*: $(\lambda x. \text{poly } p \text{ (real } x)) \sim_{[at-top]} (\lambda x. \text{lead-coeff } p * \text{real } x^{\wedge}(\text{degree } p))$
proof (cases degree p = 0)
 case False
 hence lc: lead-coeff p \neq 0 by auto
 have 1: 1 = $(\sum n \leq \text{degree } p. \text{if } n = \text{degree } p \text{ then } (1 :: \text{real}) \text{ else } 0)$ by simp
 from False show ?thesis
 proof (intro asymp-equivI', unfold poly-altdef sum-divide-distrib,
 subst 1, intro tendsto-sum, goal-cases)
 case (1 n)
 hence n = degree p \vee n < degree p by auto
 thus ?case
 proof
 assume n = degree p
 thus ?thesis using False lc
 by (simp, intro LIMSEQ-I exI[of - Suc 0], auto)
 qed (insert False lc, real-asymp)
 qed
next
 case True
 then obtain c where p: p = [:c:] by (metis degree-eq-zeroE)
 show ?thesis unfolding p by simp
qed

lemma *sum-root-unity*: fixes x :: 'a :: {comm-ring, division-ring}
 assumes $x^{\wedge}n = 1$
 shows $\text{sum } (\lambda i. x^{\wedge}i) \{..<n\} = (\text{if } x = 1 \text{ then of-nat } n \text{ else } 0)$
proof (cases x = 1 \vee n = 0)
 case x: False
 from x obtain m where n: n = Suc m by (cases n, auto)
 have id: $\{..<n\} = \{0..m\}$ unfolding n by auto
 show ?thesis using assms x n unfolding id sum-gp by (auto simp: divide-inverse)
qed auto

lemma *sum-root-unity-power-pos-implies-1*:
 assumes sumpos: $\bigwedge k. \text{Re } (\text{sum } (\lambda i. b \ i * x \ i^{\wedge}k) I) > 0$
 and root-unity: $\bigwedge i. i \in I \implies \exists d. d \neq 0 \wedge x \ i^{\wedge}d = 1$
 shows $1 \in x \text{ ' } I$
proof (rule ccontr)
 assume \neg ?thesis
 hence x: $i \in I \implies x \ i \neq 1$ for i by auto
 from sumpos[of 0] have I: finite I I \neq {}
 using sum.infinite by fastforce+
 have $\forall i. \exists d. i \in I \implies d \neq 0 \wedge x \ i^{\wedge}d = 1$ using root-unity by auto
 from choice[OF this] obtain d where d: $\bigwedge i. i \in I \implies d \ i \neq 0 \wedge x \ i^{\wedge}(d \ i)$
 = 1 by auto
 define D where D = prod d I
 have D0: 0 < D unfolding D-def


```

    by (rule prod-pos, insert d, auto)
  have  $0 < \text{sum } (\lambda k. \text{Re } (\text{sum } (\lambda i. b\ i * x\ i^{\wedge} k)\ I)) \{.. < D\}$ 
    by (rule sum-pos[OF - - sumpos], insert D0, auto)
  also have  $\dots = \text{Re } (\text{sum } (\lambda k. \text{sum } (\lambda i. b\ i * x\ i^{\wedge} k)\ I) \{.. < D\})$  by auto
  also have  $\text{sum } (\lambda k. \text{sum } (\lambda i. b\ i * x\ i^{\wedge} k)\ I) \{.. < D\}$ 
    =  $\text{sum } (\lambda i. \text{sum } (\lambda k. b\ i * x\ i^{\wedge} k) \{.. < D\})\ I$  by (rule sum.swap)
  also have  $\dots = \text{sum } (\lambda i. b\ i * \text{sum } (\lambda k. x\ i^{\wedge} k) \{.. < D\})\ I$ 
    by (rule sum.cong, auto simp: sum-distrib-left)
  also have  $\dots = 0$ 
  proof (rule sum.neutral, intro ballI)
    fix i
    assume i:  $i \in I$ 
    from d[OF this] x[OF this] have d:  $d\ i \neq 0$  and rt-unity:  $x\ i^{\wedge} d\ i = 1$ 
      and x:  $x\ i \neq 1$  by auto
    have  $\exists C. D = d\ i * C$  unfolding D-def
      by (subst prod.remove[of - i], insert i I, auto)
    then obtain C where D:  $D = d\ i * C$  by auto
    have image:  $(\bigwedge x. f\ x = x) \implies \{.. < D\} = f\ \{.. < D\}$  for f by auto
    let ?g =  $(\lambda (a,c). a + d\ i * c)$ 
    have  $\{.. < D\} = ?g\ \{.. < d\ i * C\}$ 
      unfolding image-image split D[symmetric] by (rule image, insert d mod-mult-div-eq,
blast)
    also have  $(\lambda j. (j \bmod d\ i, j \text{ div } d\ i))\ \{.. < d\ i * C\} = \{.. < d\ i\} \times \{.. < C\}$ 
      (is ?f\ \{.. < ?A\} = ?B)
    proof -
      {
        fix x
        assume x  $\in ?B$  then obtain a c where x:  $x = (a,c)$  and a:  $a < d\ i$  and
c:  $c < C$  by auto
        hence  $a + c * d\ i < d\ i * (1 + c)$  by simp
        also have  $\dots \leq d\ i * C$  by (rule mult-le-mono2, insert c, auto)
        finally have  $a + c * d\ i \in ?A$  by auto
        hence  $?f\ (a + c * d\ i) \in ?f\ \{.. < ?A\}$  by blast
        also have  $?f\ (a + c * d\ i) = x$  unfolding x using a by auto
        finally have  $x \in ?f\ \{.. < ?A\}$ .
      }
    thus ?thesis using d by (auto simp: div-lt-nat)
  qed
  finally have D:  $\{.. < D\} = (\lambda (a,c). a + d\ i * c)\ \{.. < ?B\}$  by auto
  have inj: inj-on ?g ?B
  proof -
    {
      fix a1 a2 c1 c2
      assume id:  $?g\ (a1,c1) = ?g\ (a2,c2)$  and *:  $(a1,c1) \in ?B\ (a2,c2) \in ?B$ 
      from arg-cong[OF id, of  $\lambda x. x \text{ div } d\ i$ ] * have c:  $c1 = c2$  by auto
      from arg-cong[OF id, of  $\lambda x. x \bmod d\ i$ ] * have a:  $a1 = a2$  by auto
      note a c
    }
    thus ?thesis by (smt SigmaE inj-onI)
  
```

```

qed
have sum (λ k. x i ^ k) {..< D} = sum (λ (a,c). x i ^ (a + d i * c)) ?B
  unfolding D by (subst sum.reindex, rule inj, auto intro!: sum.cong)
also have ... = sum (λ (a,c). x i ^ a) ?B
  by (rule sum.cong, auto simp: power-add power-mult rt-unity)
also have ... = 0 unfolding sum.cartesian-product[symmetric] sum.swap[of
- {..< C}]
  by (rule sum.neutral, intro ballI, subst sum-root-unity[OF rt-unity], insert x,
auto)
  finally
  show b i * sum (λ k. x i ^ k) {..< D} = 0 by simp
qed
finally show False by simp
qed

```

```

fun j-to-jb-index :: (nat × 'a)list ⇒ nat ⇒ nat × nat where
  j-to-jb-index ((n,a) # n-as) i = (if i < n then (0,i) else
    let rec = j-to-jb-index n-as (i - n) in (Suc (fst rec), snd rec))

```

```

fun jb-to-j-index :: (nat × 'a)list ⇒ nat × nat ⇒ nat where
  jb-to-j-index n-as (0,j) = j
| jb-to-j-index ((n,-) # n-as) (Suc i, j) = n + jb-to-j-index n-as (i,j)

```

```

lemma j-to-jb-index: assumes i < sum-list (map fst n-as)
  and j < sum-list (map fst n-as)
  and j-to-jb-index n-as i = (bi, li)
  and j-to-jb-index n-as j = (bj, lj)
  and n-as ! bj = (n, a)

```

```

shows ((jordan-matrix n-as) ^m r) $$ (i,j) = (if bi = bj then ((jordan-block n a)
^m r) $$ (li, lj) else 0)

```

```

  ∧ (bi = bj → li < n ∧ lj < n ∧ bj < length n-as ∧ (n,a) ∈ set n-as)

```

```

  unfolding jordan-matrix-pow using assms

```

```

proof (induct n-as arbitrary: i j bi bj)

```

```

  case (Cons mb n-as i j bi bj)

```

```

  obtain m b where mb: mb = (m,b) by force

```

```

  note Cons = Cons[unfolded mb]

```

```

  have [simp]: dim-col (case x of (n, a) ⇒ 1_m n) = fst x for x by (cases x, auto)

```

```

  have [simp]: dim-row (case x of (n, a) ⇒ 1_m n) = fst x for x by (cases x, auto)

```

```

  have [simp]: dim-col (case x of (n, a) ⇒ jordan-block n a ^m r) = fst x for x

```

```

  by (cases x, auto)

```

```

  have [simp]: dim-row (case x of (n, a) ⇒ jordan-block n a ^m r) = fst x for x

```

```

  by (cases x, auto)

```

```

  consider (UL) i < m j < m | (UR) i < m j ≥ m | (LL) i ≥ m j < m

```

```

    | (LR) i ≥ m j ≥ m by linarith

```

```

  thus ?case

```

```

proof cases

```

```

  case UL

```

```

  with Cons(2-) show ?thesis unfolding mb by (auto simp: Let-def)

```

```

next

```

```

    case UR
  with Cons(2-) show ?thesis unfolding mb by (auto simp: Let-def dim-diag-block-mat
o-def)
  next
    case LL
  with Cons(2-) show ?thesis unfolding mb by (auto simp: Let-def dim-diag-block-mat
o-def)
  next
    case LR
  let ?i = i - m
  let ?j = j - m
  from LR Cons(2-) have bi: j-to-jb-index n-as ?i = (bi - 1, li) bi ≠ 0 by
(auto simp: Let-def)
  from LR Cons(2-) have bj: j-to-jb-index n-as ?j = (bj - 1, lj) bj ≠ 0 by
(auto simp: Let-def)
  from LR Cons(2-) have i: ?i < sum-list (map fst n-as) by auto
  from LR Cons(2-) have j: ?j < sum-list (map fst n-as) by auto
  from LR Cons(2-) bj(2) have nas: n-as ! (bj - 1) = (n, a) by (cases bj,
auto)
  from bi(2) bj(2) have id: (bi - 1 = bj - 1) = (bi = bj) by auto
  note IH = Cons(1)[OF i j bi(1) bj(1) nas, unfolded id]
  have id: diag-block-mat (map (λa. case a of (n, a) ⇒ jordan-block n a ^m r)
(mb # n-as)) $$ (i, j)
    = diag-block-mat (map (λa. case a of (n, a) ⇒ jordan-block n a ^m r) n-as)
    $$ (?i, ?j)
  using i j LR unfolding mb by (auto simp: Let-def dim-diag-block-mat o-def)
  show ?thesis using IH unfolding id by auto
qed
qed auto

```

```

lemma j-to-jb-index-rev: assumes j: j-to-jb-index n-as i = (bi, li)
  and i: i < sum-list (map fst n-as)
  and k: k ≤ li
shows li ≤ i ∧ j-to-jb-index n-as (i - k) = (bi, li - k) ∧ (
j-to-jb-index n-as j = (bi, li - k) → j < sum-list (map fst n-as) → j = i - k)
  using j i
proof (induct n-as arbitrary: i bi j)
  case (Cons mb n-as i bi j)
  obtain m b where mb: mb = (m, b) by force
  note Cons = Cons[unfolded mb]
  show ?case
  proof (cases i < m)
    case True
    thus ?thesis unfolding mb using Cons(2-) by (auto simp: Let-def)
  next
    case i-large: False
    let ?i = i - m
    have i: ?i < sum-list (map fst n-as) using Cons(2-) i-large by auto
    from Cons(2-) i-large have j: j-to-jb-index n-as ?i = (bi - 1, li)

```

```

    and bi: bi ≠ 0 by (auto simp: Let-def)
  note IH = Cons(1)[OF j i]
  from IH have IH1: j-to-jb-index n-as (i - m - k) = (bi - 1, li - k) and
    li: li ≤ i - m by auto
  from li have aim1: li ≤ i by auto
  from li k i-large have i - k ≥ m by auto
  hence aim2: j-to-jb-index (mb # n-as) (i - k) = (bi, li - k)
    using IH1 bi by (auto simp: mb Let-def add commute)
  {
    assume *: j-to-jb-index (mb # n-as) j = (bi, li - k)
      j < sum-list (map fst (mb # n-as))
    from * bi have j: j ≥ m unfolding mb by (auto simp: Let-def split: if-splits)
    let ?j = j - m
    from j * have jj: ?j < sum-list (map fst n-as) unfolding mb by auto
    from j * have **: j-to-jb-index n-as (j - m) = (bi - 1, li - k) using bi mb
      by (cases j-to-jb-index n-as (j - m), auto simp: Let-def)
    from IH[of ?j] jj ** have j - m = i - m - k by auto
    with j i-large k have j = i - k using ⟨m ≤ i - k⟩ by linarith
  } note aim3 = this
  show ?thesis using aim1 aim2 aim3 by blast
qed
qed auto

```

```

locale spectral-radius-1-jnf-max =
  fixes A :: real mat and n m :: nat and lam :: complex and n-as
  assumes A: A ∈ carrier-mat n n
  and nonneg: nonneg-mat A
  and jnf: jordan-nf (map-mat complex-of-real A) n-as
  and mem: (m, lam) ∈ set n-as
  and lam1: cmod lam = 1
  and sr1:  $\bigwedge x. \text{poly} (\text{char-poly } A) x = 0 \implies x \leq 1$ 
  and max-block:  $\bigwedge k \text{ la}. (k, \text{la}) \in \text{set } n\text{-as} \implies \text{cmod } \text{la} \leq 1 \wedge (\text{cmod } \text{la} = 1 \implies k \leq m)$ 
begin

```

```

lemma n-as0: 0 ∉ fst `set n-as
  using jnf[unfolded jordan-nf-def] ..

```

```

lemma m0: m ≠ 0 using mem n-as0 by force

```

```

abbreviation cA where cA ≡ map-mat complex-of-real A
abbreviation J where J ≡ jordan-matrix n-as

```

```

lemma sim-A-J: similar-mat cA J
  using jnf[unfolded jordan-nf-def] ..

```

```

lemma sumlist-nf: sum-list (map fst n-as) = n
proof -

```

have $sum\text{-}list\ (map\ fst\ n\text{-}as) = dim\text{-}row\ (jordan\text{-}matrix\ n\text{-}as)$ **by** $simp$
also have $\dots = dim\text{-}row\ cA$ **using** $similar\text{-}matD[OF\ sim\text{-}A\text{-}J]$ **by** $auto$
finally show $?thesis$ **using** A **by** $auto$
qed

definition $p :: nat \Rightarrow real\ poly$ **where**
 $p\ s = (\prod\ i = 0..<s. [: -\ of\text{-}nat\ i / of\text{-}nat\ (s - i), 1 / of\text{-}nat\ (s - i) :])$

lemma $p\text{-}binom$: **assumes** $sk: s \leq k$
shows $of\text{-}nat\ (k\ choose\ s) = poly\ (p\ s)\ (of\text{-}nat\ k)$
unfolding $binomial\text{-}altdef\text{-}of\text{-}nat[OF\ assms]$ $p\text{-}def$ $poly\text{-}prod$
by $(rule\ prod.cong[OF\ refl], insert\ sk, auto\ simp: field\text{-}simps)$

lemma $p\text{-}binom\text{-}complex$: **assumes** $sk: s \leq k$
shows $of\text{-}nat\ (k\ choose\ s) = complex\text{-}of\text{-}real\ (poly\ (p\ s)\ (of\text{-}nat\ k))$
unfolding $p\text{-}binom[OF\ sk, symmetric]$ **by** $simp$

lemma $deg\text{-}p$: $degree\ (p\ s) = s$ **unfolding** $p\text{-}def$
by $(subst\ degree\text{-}prod\text{-}eq\text{-}sum\text{-}degree, auto)$

lemma $lead\text{-}coeff\text{-}p$: $lead\text{-}coeff\ (p\ s) = (\prod\ i = 0..<s. 1 / (of\text{-}nat\ s - of\text{-}nat\ i))$
unfolding $p\text{-}def$ $lead\text{-}coeff\text{-}prod$
by $(rule\ prod.cong[OF\ refl], auto)$

lemma $lead\text{-}coeff\text{-}p\text{-}gt\text{-}0$: $lead\text{-}coeff\ (p\ s) > 0$ **unfolding** $lead\text{-}coeff\text{-}p$
by $(rule\ prod\text{-}pos, auto)$

definition $c = lead\text{-}coeff\ (p\ (m - 1))$

lemma $c\text{-}gt\text{-}0$: $c > 0$ **unfolding** $c\text{-}def$ **by** $(rule\ lead\text{-}coeff\text{-}p\text{-}gt\text{-}0)$
lemma $c0$: $c \neq 0$ **using** $c\text{-}gt\text{-}0$ **by** $auto$

definition PP **where** $PP = (SOME\ PP. similar\text{-}mat\text{-}wit\ cA\ J\ (fst\ PP)\ (snd\ PP))$

definition P **where** $P = fst\ PP$

definition iP **where** $iP = snd\ PP$

lemma JNF : $P \in carrier\text{-}mat\ n\ n$ $iP \in carrier\text{-}mat\ n\ n$ $J \in carrier\text{-}mat\ n\ n$
 $P * iP = 1_m\ n$ $iP * P = 1_m\ n$ $cA = P * J * iP$

proof $(atomize\ (full), goal\text{-}cases)$

case 1

have $n: n = dim\text{-}row\ cA$ **using** A **by** $auto$

from $sim\text{-}A\text{-}J[unfolding\ similar\text{-}mat\text{-}def]$ **obtain** $Q\ iQ$

where $similar\text{-}mat\text{-}wit\ cA\ J\ Q\ iQ$ **by** $auto$

hence $similar\text{-}mat\text{-}wit\ cA\ J\ (fst\ (Q, iQ))\ (snd\ (Q, iQ))$ **by** $auto$

hence $similar\text{-}mat\text{-}wit\ cA\ J\ P\ iP$ **unfolding** $PP\text{-}def\ iP\text{-}def\ P\text{-}def$

by $(rule\ someI)$

from $similar\text{-}mat\text{-}witD[OF\ n\ this]$

show $?case$ **by** $auto$

qed

definition $C :: \text{nat set}$ **where**

$$C = \{j \mid j \text{ bj } lj \text{ nn } la. j < n \wedge j\text{-to-jb-index } n\text{-as } j = (bj, lj) \\ \wedge n\text{-as } ! \text{ bj} = (nn, la) \wedge cmod \text{ la} = 1 \wedge nn = m \wedge lj = nn - 1\}$$

lemma $C\text{-nonempty}$: $C \neq \{\}$

proof –

from $split\text{-list}[OF \text{ mem}]$ **obtain** $as \ bs$ **where** $n\text{-as}$: $n\text{-as} = as @ (m, lam) \# bs$
by $auto$

let $?i = sum\text{-list} (map \text{fst } as) + (m - 1)$

have $j\text{-to-jb-index } n\text{-as } ?i = (length \ as, m - 1)$

unfolding $n\text{-as}$ **by** $(induct \ as, insert \ m0, auto \ simp: Let\text{-def})$

with $lam1$ **have** $?i \in C$ **unfolding** $C\text{-def}$ **unfolding** $sumlist\text{-nf}[symmetric] \ n\text{-as}$
using $m0$ **by** $auto$

thus $?thesis$ **by** $blast$

qed

lemma $C\text{-n}$: $C \subseteq \{..<n\}$ **unfolding** $C\text{-def}$ **by** $auto$

lemma $root\text{-unity-cmod-1}$: **assumes** la : $la \in \text{snd } ' \text{ set } n\text{-as}$ **and** 1 : $cmod \ la = 1$

shows $\exists d. d \neq 0 \wedge la \wedge^d = 1$

proof –

from la **obtain** k **where** kla : $(k, la) \in \text{set } n\text{-as}$ **by** $force$

from $n\text{-as}0 \ kla$ **have** $k0$: $k \neq 0$ **by** $force$

from $split\text{-list}[OF \ kla]$ **obtain** $as \ bs$ **where** nas : $n\text{-as} = as @ (k, la) \# bs$ **by**
 $auto$

have rt : $poly (char\text{-poly } cA) \ la = 0$ **using** $k0$

unfolding $jordan\text{-nf-char-poly}[OF \ jnf]$ nas $poly\text{-prod-list } prod\text{-list-zero-iff}$ **by**
 $auto$

obtain $ks \ f$ **where** $decomp$: $decompose\text{-prod-root-unity} (char\text{-poly } A) = (ks, f)$
by $force$

from $sumlist\text{-nf}[unfolded \ nas] \ k0$ **have** $n0$: $n \neq 0$ **by** $auto$

note $pf = perron\text{-frobenius-for-complexity-jnf}(1, \gamma)[OF \ A \ n0 \ nonneg \ sr1 \ decomp,$
 $simplified]$

from $pf(1) \ pf(2)[OF \ 1 \ rt]$ **show** $\exists d. d \neq 0 \wedge la \wedge^d = 1$ **by** $metis$

qed

definition d **where** $d = (SOME \ d. \forall la. la \in \text{snd } ' \text{ set } n\text{-as} \longrightarrow cmod \ la = 1 \longrightarrow$

$$d \ la \neq 0 \wedge la \wedge^d = 1)$$

lemma d : **assumes** $(k, la) \in \text{set } n\text{-as}$ $cmod \ la = 1$

shows $la \wedge^d = 1 \wedge d \ la \neq 0$

proof –

let $?P = \lambda d. \forall la. la \in \text{snd } ' \text{ set } n\text{-as} \longrightarrow cmod \ la = 1 \longrightarrow$

$$d \ la \neq 0 \wedge la \wedge^d = 1$$

from $root\text{-unity-cmod-1}$ **have** $\forall la. \exists d. la \in \text{snd } ' \text{ set } n\text{-as} \longrightarrow cmod \ la = 1$
 \longrightarrow

$d \neq 0 \wedge \widehat{la} \ d = 1$ **by** *blast*
from *choice[OF this]* **have** $\exists d. ?P \ d$.
from *someI-ex[OF this]* **have** $?P \ d$ **unfolding** *d-def* .
from *this[rule-format, of la, OF - assms(2)] assms(1)* **show** *?thesis* **by** *force*
qed

definition *D* **where** $D = \text{prod-list } (\text{map } (\lambda \ na. \text{if } \text{cmod } (\text{snd } \ na) = 1 \text{ then } d \ (\text{snd } \ na) \text{ else } 1) \ n\text{-as})$

lemma *D0*: $D \neq 0$ **unfolding** *D-def*
by (*unfold prod-list-zero-iff, insert d, force*)

definition *f* **where** $f \ \text{off } \ k = D * k + (m-1) + \text{off}$

lemma *mono-f*: *strict-mono* (*f off*) **unfolding** *strict-mono-def f-def*
using *D0* **by** *auto*

definition *inv-op* **where** $\text{inv-op } \ \text{off } \ k = \text{inverse } (c * \text{real } (f \ \text{off } \ k) \ \widehat{(m-1)})$

lemma *limit-jordan-block*: **assumes** *kla*: $(k, la) \in \text{set } n\text{-as}$
and *ij*: $i < k \ j < k$

shows $(\lambda N. (\text{jordan-block } \ k \ la \ \widehat{m} \ (f \ \text{off } \ N))) \ \S\S \ (i, j) * \text{inv-op } \ \text{off } \ N$
 $\longrightarrow (\text{if } i = 0 \wedge j = k - 1 \wedge \text{cmod } la = 1 \wedge k = m \text{ then } la \ \widehat{\text{off}} \text{ else } 0)$

proof –

let *?c* = *of-nat* :: $\text{nat} \Rightarrow \text{complex}$
let *?r* = *of-nat* :: $\text{nat} \Rightarrow \text{real}$
let *?cr* = *complex-of-real*
from *ij* **have** *k0*: $k \neq 0$ **by** *auto*
from *jordan-nf-char-poly[OF jnf]* **have** *cA*: $\text{char-poly } \ cA = (\prod (n, a) \leftarrow n\text{-as}. [:- a, 1:] \widehat{n})$.
from *degree-monic-char-poly[OF A]* **have** $\text{degree } (\text{char-poly } \ A) = n$ **by** *auto*
have *deg*: $\text{degree } (\text{char-poly } \ cA) = n$ **using** *A* **by** (*simp add: degree-monic-char-poly*)
from *this[unfolded cA]* **have** $n = \text{degree } (\prod (n, a) \leftarrow n\text{-as}. [:- a, 1:] \widehat{n})$ **by** *auto*
also **have** $\dots = \text{sum-list } (\text{map } \text{degree } (\text{map } (\lambda(n, a). [:- a, 1:] \widehat{n}) \ n\text{-as}))$
by (*subst degree-prod-list-eq, auto*)
also **have** $\dots = \text{sum-list } (\text{map } \text{fst } \ n\text{-as})$
by (*rule arg-cong[of - - sum-list], auto simp: degree-linear-power*)
finally **have** *sum*: $\text{sum-list } (\text{map } \text{fst } \ n\text{-as}) = n$ **by** *auto*
with *split-list[OF kla]* *k0* **have** *n0*: $n \neq 0$ **by** *auto*
obtain *ks* *small* **where** *decomp*: $\text{decompose-prod-root-unity } (\text{char-poly } \ A) = (ks, \text{small})$ **by** *force*
note *pf* = *perron-frobenius-for-complexity-jnf[OF A n0 nonneg sr1 decomp]*
define *ji* **where** $ji = j - i$
have *ji*: $j - i = ji$ **unfolding** *ji-def* **by** *auto*
let *?f* = $\lambda N. c * (?r \ N) \widehat{(m-1)}$
let *?jb* = $\lambda N. (\text{jordan-block } \ k \ la \ \widehat{m} \ N) \ \S\S \ (i, j)$
let *?jbc* = $\lambda N. (\text{jordan-block } \ k \ la \ \widehat{m} \ N) \ \S\S \ (i, j) / ?f \ N$
define *e* **where** $e = (\text{if } i = 0 \wedge j = k - 1 \wedge \text{cmod } la = 1 \wedge k = m \text{ then } la \ \widehat{\text{off}} \text{ else } 0)$

```

let ?e1 =  $\lambda N :: \text{nat. } ?cr (poly (p (j - i)) (?r N)) * la \wedge (N + i - j)$ 
let ?e2 =  $\lambda N. ?cr (poly (p ji) (?r N) / ?f N) * la \wedge (N + i - j)$ 
define e2 where e2 = ?e2
let ?e3 =  $\lambda N. poly (p ji) (?r N) / (c * ?r N \wedge (m - 1)) * cmod la \wedge (N + i - j)$ 
define e3 where e3 = ?e3
define e3' where e3' =  $(\lambda N. (lead-coeff (p ji) * (?r N) \wedge ji) / (c * ?r N \wedge (m - 1))) * cmod la \wedge (N + i - j)$ 
{
  assume ij':  $i \leq j$  and la0:  $la \neq 0$ 
  {
    fix N
    assume  $N \geq k$ 
    with ij ij' have ji:  $j - i \leq N$  and id:  $N + i - j = N - ji$  unfolding ji-def
by auto
    have ?jb N =  $(?c (N \text{ choose } (j - i)) * la \wedge (N + i - j))$ 
      unfolding jordan-block-pow using ij ij' by auto
    also have ... = ?e1 N by (subst p-binom-complex[OF ji], auto)
    finally have id: ?jb N = ?e1 N .
    have ?jbc N = e2 N
    unfolding id e2-def ji-def using c-gt-0 by (simp add: norm-mult norm-divide norm-power)
  } note jbc = this
  have cmod-e2-e3:  $(\lambda n. cmod (e2 n)) \sim[at-top] e3$ 
  proof (intro asymp-equiv LIMSEQ-I exI[of - ji] allI impI)
    fix n r
    assume n:  $n \geq ji$ 
    have  $cmod (e2 n) = |poly (p ji) (?r n) / (c * ?r n \wedge (m - 1))| * cmod la \wedge (n + i - j)$ 
      unfolding e2-def norm-mult norm-power norm-of-real by simp
    also have  $|poly (p ji) (?r n) / (c * ?r n \wedge (m - 1))| = poly (p ji) (?r n) / (c * real n \wedge (m - 1))$ 
      by (intro abs-of-nonneg divide-nonneg-nonneg mult-nonneg-nonneg, insert c-gt-0, auto simp: p-binom[OF n, symmetric])
    finally have  $cmod (e2 n) = e3 n$  unfolding e3-def by auto
    thus  $r > 0 \implies norm ((if cmod (e2 n) = 0 \wedge e3 n = 0 \text{ then } 1 \text{ else } cmod (e2 n) / e3 n) - 1) < r$  by simp
  } qed
  have e3':  $e3 \sim[at-top] e3'$  unfolding e3-def e3'-def
    by (intro asymp-equiv-intros, insert poly-asymp-equiv[of p ji], unfold deg-p)
  {
    assume e3'  $\longrightarrow 0$ 
    hence e3:  $e3 \longrightarrow 0$  using e3' by (meson tendsto-asymp-equiv-cong)
    have e2  $\longrightarrow 0$ 
    by (subst tendsto-norm-zero-iff[symmetric], subst tendsto-asymp-equiv-cong[OF cmod-e2-e3], rule e3)
  } note e2-via-e3 = this

  have (e2 o f off)  $\longrightarrow e$ 

```



```

proof (cases cmod la = 1 ∧ k = m ∧ i = 0 ∧ j = k - 1)
  case False
  then consider (0) la = 0 | (small) la ≠ 0 cmod la < 1 |
    (medium) cmod la = 1 k < m ∨ i ≠ 0 ∨ j ≠ k - 1
    using max-block[OF kla] by linarith
  hence main: e2 ⟶ e
  proof cases
    case 0
    hence e0: e = 0 unfolding e-def by auto
    show ?thesis unfolding e0 0 LIMSEQ-iff e2-def ji
    proof (intro exI[of - Suc j] impI allI, goal-cases)
      case (1 r n) thus ?case by (cases n + i - j, auto)
    qed
  next
  case small
  define d where d = cmod la
  from small have d: 0 < d d < 1 unfolding d-def by auto
  have e0: e = 0 using small unfolding e-def by auto
  show ?thesis unfolding e0
  by (intro e2-via-e3, unfold e3'-def d-def[symmetric], insert d c0, real-asymp)
  next
  case medium
  with max-block[OF kla] have k ≤ m by auto
  with ij medium have ji: ji < m - 1 unfolding ji-def by linarith
  have e0: e = 0 using medium unfolding e-def by auto
  show ?thesis unfolding e0
  by (intro e2-via-e3, unfold e3'-def medium power-one mult-1-right, insert
  ji c0, real-asymp)
  qed
  show (e2 o f off) ⟶ e
  by (rule LIMSEQ-subseq-LIMSEQ[OF main mono-f])
  next
  case True
  hence large: cmod la = 1 k = m i = 0 j = k - 1 by auto
  hence e: e = la ^ off and ji: ji = m - 1 unfolding e-def ji-def by auto
  from large k0 have m0: m ≥ 1 by auto
  define m1 where m1 = m - 1
  have id: (real (m - 1) - real ia) = ?r m - 1 - ?r ia for ia using m0
  unfolding m1-def by auto
  define q where q = p m1 - monom c m1
  hence pji: p ji = q + monom c m1 unfolding q-def ji m1-def by simp
  let ?e4a = λ x. (complex-of-real (poly q (real x) / (c * real x ^ m1))) * la ^
  (x + i - j)
  let ?e4b = λ x. la ^ (x + i - j)
  {
  fix x :: nat
  assume x: x ≠ 0
  have e2 x = ?e4a x + ?e4b x
  unfolding e2-def pji poly-add poly-monom m1-def[symmetric] using c0 x

```

```

by (simp add: field-simps)
} note e2-e4 = this
have e2-e4:  $\forall_F x$  in sequentially. (e2 o f off) x = (?e4a o f off) x + (?e4b o
f off) x unfolding o-def
  by (intro eventually-sequentiallyI[of Suc 0], rule e2-e4, insert D0, auto
simp: f-def)
  have (e2 o f off)  $\longrightarrow$  0 + e
  unfolding tendsto-cong[OF e2-e4]
proof (rule tendsto-add, rule LIMSEQ-subseq-LIMSEQ[OF - mono-f])
  show ?e4a  $\longrightarrow$  0
  proof (subst tendsto-norm-zero-iff[symmetric],
  unfold norm-mult norm-power large power-one mult-1-right norm-divide
norm-of-real
  tendsto-rabs-zero-iff)
  have deg-q: degree q  $\leq$  m1 unfolding q-def using deg-p[of m1]
  by (intro degree-diff-le degree-monom-le, auto)
  have coeff-q-m1: coeff q m1 = 0 unfolding q-def c-def m1-def[symmetric]
using deg-p[of m1] by simp
  from deg-q coeff-q-m1 have deg: degree q < m1  $\vee$  q = 0 by fastforce
  have eq: ( $\lambda n$ . poly q (real n) / (c * real n ^ m1))  $\sim$ [at-top]
  ( $\lambda n$ . lead-coeff q * real n ^ degree q / (c * real n ^ m1))
  by (intro asymp-equiv-intros poly-asymp-equiv)
  show ( $\lambda n$ . poly q (?r n) / (c * ?r n ^ m1))  $\longrightarrow$  0
  unfolding tendsto-asymp-equiv-cong[OF eq] using deg
  by (standard, insert c0, real-asymp, simp)
qed
next
have id: D * x + (m - 1) + off + i - j = D * x + off for x
  unfolding ji[symmetric] ji-def using ij' by auto
from d[OF kla large(1)] have 1: la ^ d la = 1 by auto
from split-list[OF kla] obtain as bs where n-as: n-as = as @ (k,la) # bs
by auto
obtain C where D: D = d la * C unfolding D-def unfolding n-as using
large by auto
  show (?e4b o f off)  $\longrightarrow$  e
  unfolding e f-def o-def id
  unfolding power-add power-mult D 1 by auto
qed
thus ?thesis by simp
qed
also have ((e2 o f off)  $\longrightarrow$  e) = ((?jbc o f off)  $\longrightarrow$  e)
proof (rule tendsto-cong, unfold eventually-at-top-linorder, rule exI[of - k],
intro allI impI, goal-cases)
  case (1 n)
  from mono-f[of off] 1 have f off n  $\geq$  k using le-trans seq-suble by blast
  from jbc[OF this] show ?case by (simp add: o-def)
qed
finally have (?jbc o f off)  $\longrightarrow$  e .
} note part1 = this

```

```

{
  assume  $i > j \vee la = 0$ 
  hence  $e: e = 0$  and  $jbn: N \geq k \implies ?jbc\ N = 0$  for  $N$ 
  unfolding jordan-block-pow e-def using ij by auto
  have  $?jbc \longrightarrow e$  unfolding e LIMSEQ-iff by (intro exI[of - k] allI impI,
subst jbn, auto)
  from LIMSEQ-subseq-LIMSEQ[OF this mono-f]
  have  $(?jbc\ o\ f\ off) \longrightarrow e$  .
} note part2 = this
from part1 part2 have  $(?jbc\ o\ f\ off) \longrightarrow e$  by linarith
thus ?thesis unfolding e-def o-def inv-op-def by (simp add: field-simps)
qed

```

definition *lambda where lambda i = snd (n-as ! fst (j-to-jb-index n-as i))*

lemma *cmod-lambda: $i \in C \implies cmod (lambda\ i) = 1$*
 unfolding *C-def lambda-def* by *auto*

lemma *R-lambda: assumes $i: i \in C$*

shows (m, lambda i) \in set n-as

proof –

from *i[unfolded C-def]*

obtain *bi li la* where *i: $i < n$* and *jb: j -to- jb -index n -as $i = (bi, li)$*

and *nth: n -as ! $bi = (m, la)$* and *cmod $la = 1 \wedge li = m - 1$* by *auto*

hence *lam: lambda i = la* unfolding *lambda-def* by *auto*

from *j-to-jb-index[of - n-as, unfolded sumlist-nf, OF i i jb jb nth] lam*

show *?thesis* by *auto*

qed

lemma *limit-jordan-matrix: assumes $ij: i < n\ j < n$*

*shows $(\lambda N. (J \hat{\ }_m (f\ off\ N))\ \S\S (i, j) * inv-op\ off\ N)$*

$\longrightarrow (if\ j \in C \wedge i = j - (m - 1)$ then $(lambda\ j) \hat{\ }_{off}$ else $0)$

proof –

obtain *bi li* where *bi: j -to- jb -index n -as $i = (bi, li)$* by *force*

obtain *bj lj* where *bj: j -to- jb -index n -as $j = (bj, lj)$* by *force*

define *la* where *la = snd (n-as ! fst (j-to-jb-index n-as j))*

obtain *nn* where *nbj: n -as ! $bj = (nn, la)$* unfolding *la-def bj fst-conv* by (*metis* *prod.collapse*)

from *j-to-jb-index[OF ij[folded sumlist-nf] bi bj nbj]*

have *eq: $bi = bj \implies li < nn \wedge lj < nn \wedge bj < length\ n$ -as $\wedge (nn, la) \in set\ n$ -as*

and

index: $(J \hat{\ }_m\ r)\ \S\S (i, j) =$

(if $bi = bj$ then $(jordan-block\ nn\ la \hat{\ }_m\ r)\ \S\S (li, lj)$ else $0)$ for *r*

by *auto*

note *index-rev = j-to-jb-index-rev[OF bj, unfolded sumlist-nf, OF ij(2) le-refl]*

show *?thesis*

proof (*cases bi = bj*)

case *False*

hence *id: $(bi = bj) = False$* by *auto*

```

{
  assume  $j \in C$   $i = j - (m - 1)$ 
  from this[unfolded C-def]  $bj$   $n bj$  have  $i = j - lj$  by auto
  from index-rev[folded this]  $bi$  False have False by auto
}
thus ?thesis unfolding index id if-False by auto
next
  case True
  hence  $id: (bi = bj) = True$  by auto
  from eq[OF True] have  $eq: li < nn$   $lj < nn$   $(nn, la) \in set$   $n-as$   $bj < length$   $n-as$ 
by auto
  have  $(\lambda N. (J \hat{\ }_m (f \ off \ N)) \ \S\ (i, j) * inv-op \ off \ N)$ 
     $\longrightarrow$   $(if \ li = 0 \wedge lj = nn - 1 \wedge cmod \ la = 1 \wedge nn = m \ then \ la \hat{\ }_{off} \ else \ 0)$ 
    unfolding index id if-True using limit-jordan-block[OF eq(3,1,2)].
  also have  $(li = 0 \wedge lj = nn - 1 \wedge cmod \ la = 1 \wedge nn = m) = (j \in C \wedge i =$ 
 $j - (m - 1))$  (is  $?l = ?r)$ 
  proof
    assume  $?r$ 
    hence  $j \in C$  ..
    from this[unfolded C-def]  $bj$   $n bj$ 
    have  $*$ :  $nn = m \ cmod \ la = 1$   $lj = nn - 1$  by auto
    from  $\langle ?r \rangle$   $*$  have  $i = j - lj$  by auto
    with  $*$  have  $li = 0$  using index-rev bi by auto
    with  $*$  show  $?l$  by auto
  next
    assume  $?l$ 
    hence  $jI: j \in C$  using  $bj$   $n bj$   $ij$  by  $(auto \ simp: C-def)$ 
    from  $\langle ?l \rangle$  have  $li = 0$  by auto
    with index-rev[of i]  $bi$   $ij(1)$   $\langle ?l \rangle$  True
    have  $i = j - (m - 1)$  by auto
    with  $jI$  show  $?r$  by auto
  qed
  finally show ?thesis unfolding la-def lambda-def .
qed
qed

```

declare *sumlist-nf*[*simp*]

lemma *A-power-P*: $cA \hat{\ }_m k * P = P * J \hat{\ }_m k$

proof (*induct k*)

case 0

show *?case* **using** *A JNF* **by** *simp*

next

case (*Suc k*)

have $cA \hat{\ }_m \ Suc \ k * P = cA \hat{\ }_m k * cA * P$ **by** *simp*

also **have** $\dots = cA \hat{\ }_m k * (P * J * iP) * P$ **using** *JNF* **by** *simp*

also **have** $\dots = (cA \hat{\ }_m k * P) * (J * (iP * P))$

using *A JNF(1-3)* **by** (*simp add: assoc-mult-mat*[*of - n n - n - n*])

also **have** $J * (iP * P) = J$ **unfolding** *JNF* **using** *JNF* **by** *auto*

finally show *?case unfolding Suc*
using *A JNF(1-3)* **by** *(simp add: assoc-mult-mat[of - n n - n - n])*
qed

lemma *inv-op-nonneg: inv-op off k ≥ 0 unfolding inv-op-def using c-gt-0* **by** *auto*

lemma *P-nonzero-entry: assumes j: j < n*
shows $\exists i < n. P \ \$\$ (i,j) \neq 0$
proof *(rule ccontr)*
assume $\neg ?thesis$
hence $0: \bigwedge i. i < n \implies P \ \$\$ (i,j) = 0$ **by** *auto*
have $1 = (iP * P) \ \$\$ (j,j)$ **using** *j unfolding JNF* **by** *auto*
also have $\dots = (\sum i = 0..<n. iP \ \$\$ (j, i) * P \ \$\$ (i, j))$
using *j JNF(1-2)* **by** *(auto simp: scalar-prod-def)*
also have $\dots = 0$ **by** *(rule sum.neutral, insert 0, auto)*
finally show *False* **by** *auto*
qed

definition *j where j = (SOME j. j ∈ C)*

lemma *j: j ∈ C unfolding j-def using C-nonempty some-in-eq* **by** *blast*

lemma *j-n: j < n using j unfolding C-def* **by** *auto*

definition *i = (SOME i. i < n ∧ P \\$\\$ (i, j - (m - 1)) ≠ 0)*

lemma *i: i < n and P-ij0: P \\$\\$ (i, j - (m - 1)) ≠ 0*
proof $-$
from *j-n* **have** *lt: j - (m - 1) < n* **by** *auto*
show $i < n \ P \ \$\$ (i, j - (m - 1)) \neq 0$
unfolding *i-def* **using** *someI-ex[OF P-nonzero-entry[OF lt]]* **by** *auto*
qed

definition *w = P *_v unit-vec n j*

lemma *w: w ∈ carrier-vec n using JNF unfolding w-def* **by** *auto*

definition *v = map-vec cmod w*

lemma *v: v ∈ carrier-vec n unfolding v-def using w* **by** *auto*

definition *u where u = iP *_v map-vec of-real v*

lemma *u: u ∈ carrier-vec n unfolding u-def using JNF(2) v* **by** *auto*

definition *a where a j = P \\$\\$ (i, j - (m - 1)) * u \$v j* **for** *j*

lemma *main-step: 0 < Re (∑ j∈C. a j * lambda j ^ l)*

proof –

```

let ?c = complex-of-real
let ?cv = map-vec ?c
let ?cm = map-mat ?c
let ?v = ?cv v
define cc where
  cc = (λ jj. ((∑ k = 0..<n. (if k = jj - (m - 1) then P $$ (i, k) else 0)) * u
  $v jj))
  {
    fix off
    define G where G = (λ k. (A ^m f off k *v v) $v i * inv-op off k)
    define F where F = (∑ j∈C. a j * lambda j ^ off)
    {
      fix kk
      define k where k = f off kk
      have ((A ^m k) *v v) $ i * inv-op off kk = Re (?c (((A ^m k) *v v) $ i *
      inv-op off kk)) by simp
      also have ?c (((A ^m k) *v v) $ i * inv-op off kk) = ?cv ((A ^m k) *v v) $
      i * ?c (inv-op off kk)
      using i A by simp
      also have ?cv ((A ^m k) *v v) = (?cm (A ^m k) *v ?v) using A
      by (subst of-real-hom.mult-mat-vec-hom[OF - v], auto)
      also have ... = (cA ^m k *v ?v)
      by (simp add: of-real-hom.mat-hom-pow[OF A])
      also have ... = (cA ^m k *v ((P * iP) *v ?v)) unfolding JNF using v by
      auto
      also have ... = (cA ^m k *v (P *v u)) unfolding u-def
      by (subst assoc-mult-mat-vec, insert JNF v, auto)
      also have ... = (P * J ^m k *v u) unfolding A-power-P[symmetric]
      by (subst assoc-mult-mat-vec, insert u JNF(1) A, auto)
      also have ... = (P *v (J ^m k *v u))
      by (rule assoc-mult-mat-vec, insert u JNF(1) A, auto)
      finally have (A ^m k *v v) $v i * inv-op off kk = Re ((P *v (J ^m k *v u))
      $ i * inv-op off kk) by simp
      also have ... = Re (∑ jj = 0..<n.
      P $$ (i, jj) * (∑ ia = 0..<n. (J ^m k) $$ (jj, ia) * u $v ia * inv-op off
      kk))
      by (subst index-mult-mat-vec, insert JNF(1) i u, auto simp: scalar-prod-def
      sum-distrib-right[symmetric]
      mult.assoc[symmetric])
      finally have (A ^m k *v v) $v i * inv-op off kk =
      Re (∑ jj = 0..<n. P $$ (i, jj) * (∑ ia = 0..<n. (J ^m k) $$ (jj, ia) * inv-op
      off kk * u $v ia))
      unfolding k-def
      by (simp only: ac-simps)
    }
    note A-to-u = this
    have G ⟶
    Re (∑ jj = 0..<n. P $$ (i, jj) *
    (∑ ia = 0..<n. (if ia ∈ C ∧ jj = ia - (m - 1) then (lambda ia) ^ off else
  
```

```

0) * u $v ia))
  unfolding A-to-u G-def
  by (intro tendsto-intros limit-jordan-matrix, auto)
  also have (∑ jj = 0..<n. P $$ (i, jj) *
    (∑ ia = 0..<n. (if ia ∈ C ∧ jj = ia - (m - 1) then (lambda ia) ^off else
0) * u $v ia))
    = (∑ jj = 0..<n. (∑ ia ∈ C. (if ia ∈ C ∧ jj = ia - (m - 1) then P $$ (i,
jj) else 0) * ((lambda ia) ^off * u $v ia)))
  by (rule sum.cong[OF refl], unfold sum-distrib-left, subst (2) sum.mono-neutral-left[of
{0..<n}],
    insert C-n, auto intro!: sum.cong)
  also have ... = (∑ ia ∈ C. (∑ jj = 0..<n. (if jj = ia - (m - 1) then P $$
(i, jj) else 0)) * ((lambda ia) ^off * u $v ia))
  unfolding sum.swap[of - C] sum-distrib-right
  by (rule sum.cong[OF refl], auto)
  also have ... = (∑ ia ∈ C. cc ia * (lambda ia) ^off) unfolding cc-def
  by (rule sum.cong[OF refl], simp)
  also have ... = F unfolding cc-def a-def F-def
  by (rule sum.cong[OF refl], insert C-n, auto)
  finally have tend3: G → Re F .

```

```

from j j-n have jR: j ∈ C and j: j < n by auto
let ?exp = λ k. sum (λ ii. P $$ (i, ii) * (J ^m k) $$ (ii,j) {..<n})
define M where M = (λ k. cmod (?exp (f off k) * inv-op off k))
{
  fix kk
  define k where k = f off kk
  define cAk where cAk = cA ^m k
  have cAk: cAk ∈ carrier-mat n n unfolding cAk-def using A by auto
  have ((A ^m k) *v v) $ i = ((map-mat cmod cAk) *v map-vec cmod w) $ i
  unfolding v-def[symmetric] cAk-def
  by (rule arg-cong[of - - λ x. (x *v v) $ i],
    unfold of-real-hom.mat-hom-pow[OF A, symmetric],
    insert nonneg-mat-power[OF A nonneg, of k], insert i j,
    auto simp: nonneg-mat-def elements-mat-def)
  also have ... ≥ cmod ((cAk *v w) $ i)
  by (subst (1 2) index-mult-mat-vec, insert i cAk w, auto simp: scalar-prod-def
    intro!: sum-norm-le norm-mult-ineq)
  also have cAk *v w = (cAk * P) *v unit-vec n j
  unfolding w-def using JNF cAk by simp
  also have ... = P *v (J ^m k *v unit-vec n j) unfolding cAk-def A-power-P
  using JNF by (subst assoc-mult-mat-vec[of - n n - n], auto)
  also have J ^m k *v unit-vec n j = col (J ^m k) j
  by (rule eq-vecI, insert j, auto)
  also have (P *v (col (J ^m k) j)) $ i = Matrix.row P i • col (J ^m k) j
  by (subst index-mult-mat-vec, insert i JNF, auto)
  also have ... = sum (λ ii. P $$ (i, ii) * (J ^m k) $$ (ii,j) {..<n})
  unfolding scalar-prod-def by (rule sum.cong, insert i j JNF(1), auto)

```

finally have $(A \hat{\ }_m k *_v v) \$v i \geq cmod (?exp k) .$
from *mult-right-mono*[*OF this inv-op-nonneg*]
have $(A \hat{\ }_m k *_v v) \$v i * inv-op off kk \geq cmod (?exp k * inv-op off kk)$
unfolding *norm-mult*
using *inv-op-nonneg* **by** *auto*
}
hence *ge*: $(A \hat{\ }_m f off k *_v v) \$v i * inv-op off k \geq M k$ **for** *k* **unfolding** *M-def*
by *auto*
from *j* **have** *mem*: $j - (m - 1) \in \{..<n\}$ **by** *auto*
have $(\lambda k. ?exp (f off k) * inv-op off k) \longrightarrow$
 $(\sum ii < n. P \$\$ (i, ii) * (if j \in C \wedge ii = j - (m - 1) then lambda j \hat{\ } off else$
 $0))$
(is $- \longrightarrow ?sum)$
unfolding *sum-distrib-right mult.assoc*
by (*rule tendsto-sum*, *rule tendsto-mult*, *force*, *rule limit-jordan-matrix*[*OF -*
j], *auto*)
also have $?sum = P \$\$ (i, j - (m - 1)) * lambda j \hat{\ } off$
by (*subst sum.remove*[*OF - mem*], *force*, *subst sum.neutral*, *insert jR*, *auto*)
finally have *tend1*: $(\lambda k. ?exp (f off k) * inv-op off k) \longrightarrow P \$\$ (i, j - (m$
 $- 1)) * lambda j \hat{\ } off .$
have *tend2*: $M \longrightarrow cmod (P \$\$ (i, j - (m - 1)) * lambda j \hat{\ } off)$ **unfolding**
M-def
by (*rule tendsto-norm*, *rule tend1*)
define *B* **where** $B = cmod (P \$\$ (i, j - (m - 1))) / 2$
have $B: 0 < B$ **unfolding** *B-def* **using** *P-ij0* **by** *auto*
{
from *P-ij0* **have** $0: P \$\$ (i, j - (m - 1)) \neq 0$ **by** *auto*
define *E* **where** $E = cmod (P \$\$ (i, j - (m - 1)) * lambda j \hat{\ } off)$
from *cmod-lambda*[*OF jR*] 0 **have** $E: E / 2 > 0$ **unfolding** *E-def* **by** *auto*
from *tend2*[*folded E-def*] **have** *tend2*: $M \longrightarrow E .$
from *ge* **have** *ge*: $G k \geq M k$ **for** *k* **unfolding** *G-def* .
from *tend2*[*unfolded LIMSEQ-iff*, *rule-format*, *OF E*]
obtain *k'* **where** *diff*: $\bigwedge k. k \geq k' \implies norm (M k - E) < E / 2$ **by** *auto*
{
fix *k*
assume $k \geq k'$
from *diff*[*OF this*] **have** *norm*: $norm (M k - E) < E / 2 .$
have $M k \geq 0$ **unfolding** *M-def* **by** *auto*
with *E norm* **have** $M k \geq E / 2$
by (*smt real-norm-def field-sum-of-halves*)
with *ge*[*of k*] *E* **have** $G k \geq E / 2$ **by** *auto*
also have $E / 2 = B$ **unfolding** *E-def B-def j norm-mult norm-power*
cmod-lambda[*OF jR*] **by** *auto*
finally have $G k \geq B .$
}
hence $\exists k'. \forall k. k \geq k' \implies G k \geq B$ **by** *auto*
}
hence *Bound*: $\exists k'. \forall k \geq k'. B \leq G k$ **by** *auto*
from *tend3*[*unfolded LIMSEQ-iff*, *rule-format*, *of B / 2*] *B*


```

obtain  $kk$  where  $kk: \bigwedge k. k \geq kk \implies \text{norm } (G k - \text{Re } F) < B / 2$  by auto
from Bound obtain  $kk'$  where  $kk': \bigwedge k. k \geq kk' \implies B \leq G k$  by auto
define  $k$  where  $k = \max kk kk'$ 
with  $kk kk'$  have  $1: \text{norm } (G k - \text{Re } F) < B / 2 \ B \leq G k$  by auto
with  $B$  have  $\text{Re } F > 0$  by (smt real-norm-def field-sum-of-halves)
}
thus ?thesis by blast
qed

```

```

lemma main-theorem:  $(m, 1) \in \text{set } n\text{-as}$ 
proof –
from main-step have  $\text{pos}: 0 < \text{Re } (\sum_{i \in C}. a_i * \text{lambda } i^l)$  for  $l$  by auto
have  $1 \in \text{lambda } C$ 
proof (rule sum-root-unity-power-pos-implies-1[of a lambda C, OF pos])
  fix  $i$ 
  assume  $i \in C$ 
  from  $d[\text{OF } R\text{-lambda}[\text{OF } \text{this}] \text{ cmod-lambda}[\text{OF } \text{this}]]$ 
  show  $\exists d. d \neq 0 \wedge \text{lambda } i^d = 1$  by auto
qed
then obtain  $i$  where  $i: i \in C$  and  $\text{lambda } i = 1$  by auto
with  $R\text{-lambda}[\text{OF } i]$  show ?thesis by auto
qed
end

```

```

lemma nonneg-sr-1-largest-jb:
assumes nonneg: nonneg-mat A
and jnf: jordan-nf (map-mat complex-of-real A) n-as
and mem:  $(m, \text{lam}) \in \text{set } n\text{-as}$ 
and lam1: cmod lam = 1
and sr1:  $\bigwedge x. \text{poly } (\text{char-poly } A) x = 0 \implies x \leq 1$ 
shows  $\exists M. M \geq m \wedge (M, 1) \in \text{set } n\text{-as}$ 
proof –
note  $\text{jnf}' = \text{jnf}[\text{unfolded } \text{jordan-nf-def}]$ 
from  $\text{jnf}'$  similar-matD[OF jnf'[THEN conjunct2]] obtain  $n$ 
where  $A: A \in \text{carrier-mat } n \ n$  and  $n\text{-as0}: 0 \notin \text{fst } n\text{-as}$  by auto
let  $?M = \{ m. \exists \text{lam}. (m, \text{lam}) \in \text{set } n\text{-as} \wedge \text{cmod } \text{lam} = 1 \}$ 
have  $m: m \in ?M$  using mem lam1 by auto
have fin: finite ?M
by (rule finite-subset[OF - finite-set[of map fst n-as]], force)
define  $M$  where  $M = \text{Max } ?M$ 
have  $M \in ?M$  using fin m unfolding M-def using Max-in by blast
then obtain  $\text{lambda}$  where  $M: (M, \text{lambda}) \in \text{set } n\text{-as} \ \text{cmod } \text{lambda} = 1$  by
auto
from  $m$  fin have  $mM: m \leq M$  unfolding M-def by simp
interpret spectral-radius-1-jnf-max A n M lambda
proof (unfold-locals, rule A, rule nonneg, rule jnf, rule M, rule M, rule sr1)
  fix  $k \ \text{la}$ 
  assume  $\text{kla}: (k, \text{la}) \in \text{set } n\text{-as}$ 

```

with *fin* **have** $1: \text{cmod } la = 1 \longrightarrow k \leq M$ **unfolding** *M-def* **using** *Max-ge* **by**
blast
obtain *ks f* **where** *decomp: decompose-prod-root-unity (char-poly A) = (ks, f)*
by *force*
from *n-as0 kla* **have** $k0: k \neq 0$ **by** *force*
let $?cA = \text{map-mat } \text{complex-of-real } A$
from *split-list[OF kla]* **obtain** *as bs* **where** $n\text{-as} = \text{as } @ (k, la) \# bs$ **by**
auto
have $rt: \text{poly } (\text{char-poly } ?cA) la = 0$ **using** *k0*
unfolding *jordan-nf-char-poly[OF jnf] nas poly-prod-list prod-list-zero-iff* **by**
auto
have $\text{sumlist-nf}: \text{sum-list } (\text{map } \text{fst } n\text{-as}) = n$
proof –
have $\text{sum-list } (\text{map } \text{fst } n\text{-as}) = \text{dim-row } (\text{jordan-matrix } n\text{-as})$ **by** *simp*
also have $\dots = \text{dim-row } ?cA$ **using** *similar-matD[OF jnf'[THEN conjunct2]]*
by *auto*
finally show *?thesis* **using** *A* **by** *auto*
qed
from *this[unfolded nas] k0* **have** $n0: n \neq 0$ **by** *auto*
from *perron-frobenius-for-complexity-jnf(4)[OF A n0 nonneg sr1 decomp rt]*
have $\text{cmod } la \leq 1$.
with *1* **show** $\text{cmod } la \leq 1 \wedge (\text{cmod } la = 1 \longrightarrow k \leq M)$ **by** *auto*
qed
from *main-theorem*
show *?thesis* **using** *mM* **by** *auto*
qed
hide-const(open) *spectral-radius*

lemma (*in ring-hom*) *hom-smult-mat: mat_h (a ·_m A) = hom a ·_m mat_h A*
by (*rule eq-matI, auto simp: hom-mult*)

lemma *root-char-poly-smult: fixes A :: complex mat*
assumes *A: A ∈ carrier-mat n n*
and *k: k ≠ 0*
shows $(\text{poly } (\text{char-poly } (k \cdot_m A)) x = 0) = (\text{poly } (\text{char-poly } A) (x / k) = 0)$
using *order-char-poly-smult[OF A k, of x]*
by (*metis A degree-0 degree-monic-char-poly monic-degree-0 order-root smult-carrier-mat*)

theorem *real-nonneg-mat-spectral-radius-largest-jordan-block:*
assumes *real-nonneg-mat A*
and *jordan-nf A n-as*
and $(m, lam) \in \text{set } n\text{-as}$
and $\text{cmod } lam = \text{spectral-radius } A$
shows $\exists M \geq m. (M, \text{of-real } (\text{spectral-radius } A)) \in \text{set } n\text{-as}$
proof –
from *similar-matD[OF assms(2)[unfolded jordan-nf-def, THEN conjunct2]]* **ob-**
tain *n* **where**
A: A ∈ carrier-mat n n **by** *auto*
let $?c = \text{complex-of-real}$

```

define B where B = map-mat Re A
have B: B ∈ carrier-mat n n unfolding B-def using A by auto
have AB: A = map-mat ?c B unfolding B-def using assms(1)
  by (auto simp: real-nonneg-mat-def elements-mat-def)
have nonneg: nonneg-mat B using assms(1) unfolding AB
  by (auto simp: real-nonneg-mat-def elements-mat-def nonneg-mat-def)
let ?sr = spectral-radius A
show ?thesis
proof (cases ?sr = 0)
  case False
  define isr where isr = inverse ?sr
  let ?nas = map (λ(n, a). (n, ?c isr * a)) n-as
  from False have isr0: isr ≠ 0 unfolding isr-def by auto
  hence cisr0: ?c isr ≠ 0 by auto
  from False assms(4) have isr-pos: isr > 0 unfolding isr-def
    by (smt norm-ge-zero positive-imp-inverse-positive)
  define C where C = isr ·m B
  have C: C ∈ carrier-mat n n using B unfolding C-def by auto
  have BC: B = ?sr ·m C using isr0 unfolding C-def isr-def by auto
  have nonneg: nonneg-mat C unfolding C-def using isr-pos nonneg
    unfolding nonneg-mat-def elements-mat-def by auto
  from jordan-nf-smult[OF assms(2)[unfolded AB] cisr0]
  have jnf: jordan-nf (map-mat ?c C) ?nas unfolding C-def by (auto simp:
of-real-hom.hom-smult-mat)
  from assms(3) have mem: (m, ?c isr * lam) ∈ set ?nas by auto
  have 1: cmod (?c isr * lam) = 1 using False isr-pos unfolding isr-def
norm-mult assms(4)
  by (smt mult.commute norm-of-real right-inverse)
  {
  fix x
  have B': map-mat ?c B ∈ carrier-mat n n using B by auto
  assume poly (char-poly C) x = 0
  hence poly (char-poly (map-mat ?c C)) (?c x) = 0 unfolding of-real-hom.char-poly-hom[OF
C] by auto
  hence poly (char-poly A) (?c x / ?c isr) = 0 unfolding C-def of-real-hom.hom-smult-mat
AB
  unfolding root-char-poly-smult[OF B' cisr0] .
  hence eigenvalue A (?c x / ?c isr) unfolding eigenvalue-root-char-poly[OF
A] .
  hence mem: cmod (?c x / ?c isr) ∈ cmod ' spectrum A unfolding spectrum-def
by auto
  from Max-ge[OF finite-imageI this]
  have cmod (?c x / ?c isr) ≤ ?sr unfolding Spectral-Radius.spectral-radius-def
  using A card-finite-spectrum(1) by blast
  hence cmod (?c x) ≤ 1 using isr0 isr-pos unfolding isr-def
  by (auto simp: field-simps norm-divide norm-mult)
  hence x ≤ 1 by auto
  } note sr = this
from nonneg-sr-1-largest-jb[OF nonneg jnf mem 1 sr] obtain M where

```

```

    M:  $M \geq m$  ( $M, 1$ )  $\in$  set ?nas by blast
    from  $M(2)$  obtain a where mem:  $(M, a) \in$  set n-as and  $1 = ?c$  isr * a by
    auto
    from this(2) have a:  $a = ?c$  ?sr using isr0 unfolding isr-def by (auto simp:
    field-simps)
    show ?thesis
    by (intro exI[of - M], insert mem a M(1), auto)
  next
  case True
  from jordan-nf-root-char-poly[OF assms(2,3)]
  have eigenvalue A lam unfolding eigenvalue-root-char-poly[OF A] .
  hence cmod lam  $\in$  cmod ' spectrum A unfolding spectrum-def by auto
  from Max-ge[OF finite-imageI this]
  have cmod lam  $\leq$  ?sr unfolding Spectral-Radius.spectral-radius-def
  using A card-finite-spectrum(1) by blast
  from this[unfolded True] have lam0: lam = 0 by auto
  show ?thesis unfolding True using assms(3)[unfolded lam0] by auto
qed
qed
end

```

8 Homomorphisms of Gauss-Jordan Elimination, Kernel and More

```

theory Hom-Gauss-Jordan
  imports Jordan-Normal-Form.Matrix-Kernel
    Jordan-Normal-Form.Jordan-Normal-Form-Uniqueness
begin

lemma (in comm-ring-hom) similar-mat-wit-hom: assumes
  similar-mat-wit A B C D
shows similar-mat-wit (math A) (math B) (math C) (math D)
proof -
  obtain n where n:  $n = \dim\text{-row } A$  by auto
  note * = similar-mat-witD[OF n assms]
  from * have [simp]:  $\dim\text{-row } C = n$  by auto
  note C = *(6) note D = *(7)
  note id = mat-hom-mult[OF C D] mat-hom-mult[OF D C]
  note ** = *(1-3)[THEN arg-cong[of - - math], unfolded id]
  note mult = mult-carrier-mat[of - n n]
  note hom-mult = mat-hom-mult[of - n n - n]
  show ?thesis unfolding similar-mat-wit-def Let-def unfolding ***(3) using
  ***(1,2)
  by (auto simp: n[symmetric] hom-mult simp: *(4-) mult)
qed

```

```

lemma (in comm-ring-hom) similar-mat-hom:

```

```

similar-mat A B  $\implies$  similar-mat (math A) (math B)
using similar-mat-wit-hom[of A B C D for C D]
by (smt similar-mat-def)

context field-hom
begin
lemma hom-swaprows: i < dim-row A  $\implies$  j < dim-row A  $\implies$ 
  swaprows i j (math A) = math (swaprows i j A)
unfolding mat-swaprows-def by (rule eq-matI, auto)

lemma hom-gauss-jordan-main: A  $\in$  carrier-mat nr nc  $\implies$  B  $\in$  carrier-mat nr
nc2  $\implies$ 
  gauss-jordan-main (math A) (math B) i j =
  map-prod math math (gauss-jordan-main A B i j)
proof (induct A B i j rule: gauss-jordan-main.induct)
  case (1 A B i j)
  note IH = 1(1-4)
  note AB = 1(5-6)
  from AB have dim: dim-row A = nr dim-col A = nc by auto
  let ?h = math
  let ?hp = map-prod math math
  show ?case unfolding gauss-jordan-main.simps[of A B i j] gauss-jordan-main.simps[of
  ?h A - i j]
    index-map-mat Let-def if-distrib[of ?hp] dim
  proof (rule if-cong[OF refl], goal-cases)
    case 1
    note IH = IH[OF dim[symmetric] 1 refl]
    from 1 have ij: i < nr j < nc by auto
    hence hij: (?h A) $$ (i,j) = hom (A $$ (i,j)) using AB by auto
    define ixs where ixs = concat (map (λi'. if A $$ (i', j) ≠ 0 then [i'] else []))
    [Suc i..nr]
    have id: map (λi'. if math A $$ (i', j) ≠ 0 then [i'] else []) [Suc i..nr] =
    map (λi'. if A $$ (i', j) ≠ 0 then [i'] else []) [Suc i..nr]
    by (rule map-cong[OF refl], insert ij AB, auto)
    show ?case unfolding hij hom-0-iff hom-1-iff id ixs-def[symmetric]
    proof (rule if-cong[OF refl - if-cong[OF refl]], goal-cases)
      case 1
      note IH = IH(1,2)[OF 1, folded ixs-def]
      show ?case
      proof (cases ixs)
        case Nil
        show ?thesis unfolding Nil using IH(1)[OF Nil AB] by auto
      next
      case (Cons I ix)
      hence I  $\in$  set ixs by auto
      hence I: I < nr unfolding ixs-def by auto
      from AB have swap: swaprows i I A  $\in$  carrier-mat nr nc swaprows i I B  $\in$ 
carrier-mat nr nc2
      by auto

```

```

      show ?thesis unfolding Cons list.simps IH(2)[OF Cons swap,symmetric]
using AB ij I
  by (auto simp: hom-swaprows)
qed
next
case 2
from AB have elim: eliminate-entries ( $\lambda i. A \ \$\$ (i, j)$ )  $A \ i \ j \in \text{carrier-mat}$ 
nr nc
  eliminate-entries ( $\lambda i. A \ \$\$ (i, j)$ )  $B \ i \ j \in \text{carrier-mat nr nc2}$ 
  unfolding eliminate-entries-gen-def by auto
  show ?case unfolding IH(3)[OF 2 refl elim, symmetric]
  by (rule arg-cong2[of - - -  $\lambda x y. \text{gauss-jordan-main } x \ y \ (\text{Suc } i) \ (\text{Suc } j)$ ];
    intro eq-matI, insert AB ij, auto simp: eliminate-entries-gen-def hom-minus
hom-mult)
  next
  case 3
  from AB have mult: multrow i (inverse (A  $\ \$\$ (i, j)$ ))  $A \in \text{carrier-mat nr nc}$ 
    multrow i (inverse (A  $\ \$\$ (i, j)$ ))  $B \in \text{carrier-mat nr nc2}$  by auto
  show ?case unfolding IH(4)[OF 3 refl mult, symmetric]
  by (rule arg-cong2[of - - -  $\lambda x y. \text{gauss-jordan-main } x \ y \ i \ j$ ];
    intro eq-matI, insert AB ij, auto simp: hom-inverse hom-mult)
  qed
qed auto
qed

```

lemma *hom-gauss-jordan*: $A \in \text{carrier-mat nr nc} \implies B \in \text{carrier-mat nr nc2} \implies$
 $\text{gauss-jordan } (\text{mat}_h A) \ (\text{mat}_h B) = \text{map-prod } \text{mat}_h \ \text{mat}_h \ (\text{gauss-jordan } A \ B)$
unfolding *gauss-jordan-def* **using** *hom-gauss-jordan-main* **by** *blast*

lemma *hom-gauss-jordan-single[simp]*: $\text{gauss-jordan-single } (\text{mat}_h A) = \text{mat}_h \ (\text{gauss-jordan-single } A)$

proof –
let $?nr = \text{dim-row } A$ **let** $?nc = \text{dim-col } A$
have $0: 0_m \ ?nr \ 0 \in \text{carrier-mat } ?nr \ 0$ **by** *auto*
have $\text{dim}: \text{dim-row } (\text{mat}_h A) = ?nr$ **by** *auto*
have $\text{hom}0: \text{mat}_h (0_m \ ?nr \ 0) = 0_m \ ?nr \ 0$ **by** *auto*
have $A: A \in \text{carrier-mat } ?nr \ ?nc$ **by** *auto*
from *hom-gauss-jordan*[*OF A 0*] A
show ?thesis **unfolding** *gauss-jordan-single-def dim hom0* **by** (*metis fst-map-prod*)
qed

lemma *hom-pivot-positions-main-gen*: **assumes** $A: A \in \text{carrier-mat nr nc}$
shows *pivot-positions-main-gen 0* ($\text{mat}_h A$) $nr \ nc \ i \ j = \text{pivot-positions-main-gen}$
 $0 \ A \ nr \ nc \ i \ j$

proof (*induct rule: pivot-positions-main-gen.induct*[*of nr nc A 0*])
case ($1 \ i \ j$)
note $IH = \text{this}$
show ?case **unfolding** *pivot-positions-main-gen.simps*[*of - - nr nc i j*]
proof (*rule if-cong*[*OF refl if-cong*[*OF refl - refl*] *refl*], *goal-cases*)

```

    case 1
    with A have id: (math A) $$ (i,j) = hom (A $$ (i,j)) by simp
    note IH = IH[OF 1]
    show ?case unfolding id hom-0-iff
      by (rule if-cong[OF refl IH(1)], force, subst IH(2), auto)
  qed
qed

lemma hom-pivot-positions[simp]: pivot-positions (math A) = pivot-positions A
  unfolding pivot-positions-def by (subst hom-pivot-positions-main-gen, auto)

lemma hom-kernel-dim[simp]: kernel-dim (math A) = kernel-dim A
  unfolding kernel-dim-code by simp

lemma hom-char-matrix: assumes A: A ∈ carrier-mat n n
  shows char-matrix (math A) (hom x) = math (char-matrix A x)
  unfolding char-matrix-def
  by (rule eq-matI, insert A, auto simp: hom-minus)

lemma hom-dim-gen-eigenspace: assumes A: A ∈ carrier-mat n n
  shows dim-gen-eigenspace (math A) (hom x) = dim-gen-eigenspace A x
proof (intro ext)
  fix k
  show dim-gen-eigenspace (math A) (hom x) k = dim-gen-eigenspace A x k
    unfolding dim-gen-eigenspace-def hom-char-matrix[OF A]
      mat-hom-pow[OF char-matrix-closed[OF A], symmetric] by simp
qed
end
end

```

9 Combining Spectral Radius Theory with Perron Frobenius theorem

```

theory Spectral-Radius-Theory-2
imports
  Spectral-Radius-Largest-Jordan-Block
  Hom-Gauss-Jordan
begin

```

```

hide-const(open) Coset.order

```

```

lemma jnf-complexity-generic: fixes A :: complex mat
  assumes A: A ∈ carrier-mat n n
  and sr:  $\bigwedge x. \text{poly} (\text{char-poly } A) x = 0 \implies \text{cmod } x \leq 1$ 
  and 1:  $\bigwedge x. \text{poly} (\text{char-poly } A) x = 0 \implies \text{cmod } x = 1 \implies$ 
    order x (char-poly A) > d + 1  $\implies$ 
    ( $\forall \text{ bsize} \in \text{fst 'set (compute-set-of-jordan-blocks } A \ x). \text{ bsize} \leq d + 1$ )
  shows  $\exists c1 \ c2. \forall k. \text{norm-bound } (A \hat{=} \text{ }_m k) (c1 + c2 * \text{of-nat } k \hat{=} d)$ 

```

```

proof -
  from char-poly-factorized[OF A] obtain as where cA: char-poly A = ( $\prod$  a $\leftarrow$ as.
[: - a, 1:])
  and lenn: length as = n by auto
  from jordan-nf-exists[OF A cA] obtain n-xs where jnf: jordan-nf A n-xs ..
  have dd:  $x^d = x^{(d+1)-1}$  for x by simp
  show ?thesis unfolding dd
  proof (rule jordan-nf-matrix-poly-bound[OF A - - jnf])
    fix n x
    assume nx: (n,x)  $\in$  set n-xs
    from jordan-nf-block-size-order-bound[OF jnf nx]
    have no:  $n \leq$  order x (char-poly A) by auto
    {
      assume 0 < n
      with no have order x (char-poly A)  $\neq$  0 by auto
      hence rt: poly (char-poly A) x = 0 unfolding order-root by auto
      from sr[OF this] show cmod x  $\leq$  1 .
      note rt
    } note sr = this
    assume c1: cmod x = 1
    show  $n \leq d + 1$ 
    proof (rule ccontr)
      assume  $\neg n \leq d + 1$ 
      hence lt:  $n > d + 1$  by auto
      with sr have rt: poly (char-poly A) x = 0 by auto
      from lt no have ord:  $d + 1 <$  order x (char-poly A) by auto
      from 1[OF rt c1 ord, unfolded compute-set-of-jordan-blocks[OF jnf]] nx lt
      show False by force
    qed
  qed
qed

```

```

lemma norm-bound-complex-to-real: fixes A :: real mat
  assumes A: A  $\in$  carrier-mat n n
  and bnd:  $\exists$  c1 c2.  $\forall$  k. norm-bound ((map-mat complex-of-real A)  $\hat{\ }_m$  k) (c1 +
c2 * of-nat k  $\hat{\ }^d$ )
  shows  $\exists$  c1 c2.  $\forall$  k a. a  $\in$  elements-mat (A  $\hat{\ }_m$  k)  $\longrightarrow$  abs a  $\leq$  (c1 + c2 * of-nat
k  $\hat{\ }^d$ )
  proof -
    let ?B = map-mat complex-of-real A
    from bnd obtain c1 c2 where nb:  $\bigwedge$  k. norm-bound (?B  $\hat{\ }_m$  k) (c1 + c2 * real
k  $\hat{\ }^d$ ) by auto
    show ?thesis
    proof (rule exI[of - c1], rule exI[of - c2], intro allI impI)
      fix k a
      assume a  $\in$  elements-mat (A  $\hat{\ }_m$  k)
      with pow-carrier-mat[OF A] obtain i j where a: a = (A  $\hat{\ }_m$  k) $$ (i,j) and
ij:  $i < n$   $j < n$ 
      unfolding elements-mat by force
    qed
  qed

```



```

from ij nb[of k] A have norm ((?B  $\hat{m}$  k)  $\S\S$  (i,j))  $\leq c1 + c2 * \text{real } k \hat{d}$ 
  unfolding norm-bound-def by auto
also have (?B  $\hat{m}$  k)  $\S\S$  (i,j) = of-real a
  unfolding of-real-hom.mat-hom-pow[OF A, symmetric] a using ij A by auto
also have norm (complex-of-real a) = abs a by auto
finally show abs a  $\leq (c1 + c2 * \text{real } k \hat{d})$ .
qed
qed

lemma dim-gen-eigenspace-max-jordan-block: assumes jnf: jordan-nf A n-as
shows dim-gen-eigenspace A l d = order l (char-poly A)  $\longleftrightarrow$ 
  ( $\forall n. (n,l) \in \text{set } n\text{-as} \longrightarrow n \leq d$ )
proof -
let ?list = [(n, e) $\leftarrow$ n-as . e = l]
define list where list = [na $\leftarrow$ n-as . snd na = l]
have list: ?list = list unfolding list-def by (induct n-as, force+)
have id: ( $\forall n. (n, l) \in \text{set } n\text{-as} \longrightarrow n \leq d$ ) = ( $\forall n \in \text{set } (\text{map } \text{fst } \text{list}). n \leq d$ )
  unfolding list-def by auto
define ns where ns = map fst list
show ?thesis
  unfolding dim-gen-eigenspace[OF jnf] jordan-nf-order[OF jnf] list list-def[symmetric]
id
  unfolding ns-def[symmetric]
proof (induct ns)
case (Cons n ns)
show ?case
proof (cases n  $\leq$  d)
  case True
  thus ?thesis using Cons by auto
next
  case False
  hence n > d by auto
  moreover have sum-list (map (min d) ns)  $\leq$  sum-list ns by (induct ns, auto)
  ultimately show ?thesis by auto
qed
qed auto
qed

lemma jnf-complexity-1-complex: fixes A :: complex mat
assumes A: A  $\in$  carrier-mat n n
and nonneg: real-nonneg-mat A
and sr:  $\bigwedge x. \text{poly } (\text{char-poly } A) x = 0 \implies \text{cmod } x \leq 1$ 
and 1:  $\text{poly } (\text{char-poly } A) 1 = 0 \implies$ 
  order 1 (char-poly A) > d + 1  $\implies$ 
  dim-gen-eigenspace A 1 (d+1) = order 1 (char-poly A)
shows  $\exists c1 c2. \forall k. \text{norm-bound } (A \hat{m} k) (c1 + c2 * \text{of-nat } k \hat{d})$ 
proof -
from char-poly-factorized[OF A] obtain as where cA: char-poly A = ( $\prod a \leftarrow as.$ 
  [: $-$  a, 1:])
```

```

    and lenn: length as = n by auto
  from jordan-nf-exists[OF A cA] obtain n-as where jnf: jordan-nf A n-as ..
  have dd:  $x^d = x^{(d+1)-1}$  for x by simp
  let ?n = n
  show ?thesis unfolding dd
  proof (rule jordan-nf-matrix-poly-bound[OF A - - jnf])
    fix n a
    assume na:  $(n, a) \in \text{set } n\text{-as}$ 
    from jordan-nf-root-char-poly[OF jnf na]
    have rt: poly (char-poly A) a = 0 by auto
    with degree-monic-char-poly[OF A] have n0: ?n > 0
      by (cases ?n, auto dest: degree0-coeffs)
    from sr[OF rt] show cmod a ≤ 1 .
    assume a: cmod a = 1
    from rt have a ∈ spectrum A using A spectrum-root-char-poly by auto
    hence 11: 1 ∈ cmod ' spectrum A using a by auto
    note spec = spectral-radius-mem-max[OF A n0]
    from spec(2)[OF 11] have le: 1 ≤ spectral-radius A .
    from spec(1)[unfolded spectrum-root-char-poly[OF A]] sr have spectral-radius
    A ≤ 1 by auto
    with le have sr: spectral-radius A = 1 by auto
    show n ≤ d + 1
    proof (rule ccontr)
      assume ¬ ?thesis
      hence nd: n > d + 1 by auto
      from real-nonneg-mat-spectral-radius-largest-jordan-block[OF nonneg jnf na,
      unfolded sr a]
      obtain N where N: N ≥ n and mem:  $(N, 1) \in \text{set } n\text{-as}$  by auto
      from jordan-nf-root-char-poly[OF jnf mem] have rt: poly (char-poly A) 1 =
      0 .
      from jordan-nf-block-size-order-bound[OF jnf mem] have N ≤ order 1
      (char-poly A) .
      with N nd have d + 1 < order 1 (char-poly A) by simp
      from 1[OF rt this, unfolded dim-gen-eigenspace-max-jordan-block[OF jnf]]
      mem N nd
      show False by force
    qed
  qed
  qed

```

```

lemma jnf-complexity-1-real: fixes A :: real mat
  assumes A: A ∈ carrier-mat n n
  and nonneg: nonneg-mat A
  and sr:  $\bigwedge x. \text{poly } (\text{char-poly } A) x = 0 \implies x \leq 1$ 
  and jb: poly (char-poly A) 1 = 0  $\implies$ 
    order 1 (char-poly A) > d + 1  $\implies$ 
    dim-gen-eigenspace A 1 (d+1) = order 1 (char-poly A)
  shows  $\exists c1 c2. \forall k a. a \in \text{elements-mat } (A \hat{=} k) \implies |a| \leq c1 + c2 * \text{real } k^d$ 
  proof -

```

```

let ?c = complex-of-real
let ?A = map-mat ?c A
have A': ?A ∈ carrier-mat n n using A by auto
have nn: real-nonneg-mat ?A using nonneg A unfolding nonneg-mat-def real-nonneg-mat-def

  by (force simp: elements-mat)
have 1: 1 = ?c 1 by auto
note cp = of-real-hom.char-poly-hom[OF A]
have hom: map-poly-inj-idom-divide-hom complex-of-real ..
show ?thesis
proof (rule norm-bound-complex-to-real[OF A jnf-complexity-1-complex[OF A'
nn]],
  unfold cp of-real-hom.poly-map-poly-1, unfold 1
  of-real-hom.hom-dim-gen-eigenspace[OF A]
  map-poly-inj-idom-divide-hom.order-hom[OF hom], goal-cases)
  case 2
  thus ?case using jb by auto
next
  case (1 x)
  let ?cp = char-poly A
  assume rt: poly (map-poly ?c ?cp) x = 0
  with degree-monic-char-poly[OF A', unfolded cp] have n0: n ≠ 0
    using degree0-coeffs[of ?cp] by (cases n, auto)
  from perron-frobenius-nonneg[OF A nonneg n0]
  obtain sr ks f where sr0: 0 ≤ sr and ks: 0 ∉ set ks ks ≠ []
    and cp: ?cp = (∏ k←ks. monom 1 k - [:sr ^ k:]) * f
    and rtf: poly (map-poly ?c f) x = 0 ⇒ cmod x < sr by auto
  have sr-rt: poly ?cp sr = 0 unfolding cp poly-prod-list-zero-iff poly-mult-zero-iff
using ks
  by (cases ks, auto simp: poly-monom)
  from sr[OF sr-rt] have sr1: sr ≤ 1 .
  interpret c: map-poly-comm-ring-hom ?c ..
  from rt[unfolded cp c.hom-mult c.hom-prod-list poly-mult-zero-iff poly-prod-list-zero-iff]

show cmod x ≤ 1
proof (standard, goal-cases)
  case 2
  with rtf sr1 show ?thesis by auto
next
  case 1
  from this ks obtain p where p: p ∈ set ks
    and rt: poly (map-poly ?c (monom 1 p - [:sr ^ p:])) x = 0 by auto
  from p ks(1) have p: p ≠ 0 by metis
  from rt have x^p = (?c sr) ^ p unfolding c.hom-minus
    by (simp add: poly-monom of-real-hom.map-poly-pCons-hom)
  hence cmod x = cmod (?c sr) using p power-eq-imp-eq-norm by blast
  with sr0 sr1 show cmod x ≤ 1 by auto
qed
qed

```

qed
end

10 An efficient algorithm to compute the growth rate of A^n .

theory *Check-Matrix-Growth*

imports

Spectral-Radius-Theory-2

Sturm-Sequences.Sturm-Method

begin

hide-const (**open**) *Coset.order*

definition *check-matrix-complexity* :: *real mat* \Rightarrow *real poly* \Rightarrow *nat* \Rightarrow *bool* **where**
check-matrix-complexity *A cp d* = (*count-roots-above cp 1* = 0
 \wedge (*poly cp 1* = 0 \longrightarrow (*let ord = order 1 cp in*
 $d + 1 < \text{ord} \longrightarrow \text{kernel-dim } ((A - 1_m (\text{dim-row } A)) \hat{\ }_m (d + 1)) = \text{ord}$)))

lemma *check-matrix-complexity*: **assumes** *A*: *A* \in *carrier-mat n n* **and** *nn*: *non-neg-mat A*

and *check*: *check-matrix-complexity A (char-poly A) d*

shows $\exists c1\ c2. \forall k\ a. a \in \text{elements-mat } (A \hat{\ }_m k) \longrightarrow \text{abs } a \leq (c1 + c2 * \text{of-nat } k \hat{\ }^d)$

proof (*rule jnf-complexity-1-real[OF A nn]*)

have *id*: *dim-gen-eigenspace A 1 (d + 1) = kernel-dim ((A - 1_m (dim-row A)) $\hat{\ }_m (d + 1))$*

unfolding *dim-gen-eigenspace-def*

by (*rule arg-cong[of - - $\lambda x. \text{kernel-dim } (x \hat{\ }_m (d + 1))$], *unfold char-matrix-def*, *insert A, auto*)*

note *check* = *check[unfolded check-matrix-complexity-def*

Let-def count-roots-above-correct, folded id]

have *fin*: *finite {x. poly (char-poly A) x = 0}*

by (*rule poly-roots-finite, insert degree-monic-char-poly[OF A], auto*)

from *check* **have** *card {x. 1 < x \wedge poly (char-poly A) x = 0} = 0* **by** *auto*

from *this[unfolded card-eq-0-iff] fin*

have *{x. 1 < x \wedge poly (char-poly A) x = 0} = {}* **by** *auto*

thus *poly (char-poly A) x = 0 \implies x \leq 1* **for** *x* **by** *force*

assume *poly (char-poly A) 1 = 0* $d + 1 < \text{order 1 (char-poly A)}$

with *check* **show** *dim-gen-eigenspace A 1 (d + 1) = order 1 (char-poly A)* **by** *auto*

qed

end

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