

Properties of Orderings and Lattices

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Abstract

These components add further fundamental order and lattice-theoretic concepts and properties to Isabelle’s libraries. They follow by and large the introductory sections of the *Compendium of Continuous Lattices*, covering directed and filtered sets, down-closed and up-closed sets, ideals and filters, Galois connections, closure and co-closure operators. Some emphasis is on duality and morphisms between structures—as in the Compendium. To this end, three ad-hoc approaches to duality are compared.

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1 Introductory Remarks

Basic order- and lattice-theoretic concepts are well covered in Isabelle’s libraries, and widely used. More advanced components are spread out over various sites (e.g. [11, 9, 8, 1, 4, 2]).

This formalisation takes the initial steps towards a modern structural approach to orderings and lattices, as for instance in denotational semantics of programs, algebraic logic or pointfree topology. Building on the components for orderings and lattices in Isabelle’s main libraries, it follows the classical textbook *A Compendium of Continuous Lattices* [3] and, to a lesser extent, Johnstone’s monograph on *Stone Spaces* [5]. By integrating material from other sources and extending it, a formalisation of undergraduate-level textbook material on orderings and lattices might eventually emerge.

In the textbooks mentioned, concepts such as dualities, isomorphisms between structures and relationships between categories are emphasised. These are essential to modern mathematics beyond orderings and lattices; their formalisation with interactive theorem provers is therefore of wider interest. Nevertheless such notions seem rather underexplored with Isabelle, and I am not aware of a standard way of modelling and using them. The present setting is perhaps the simplest one in which their formalisation can be studied.

These components use Isabelle’s axiomatic approach without carrier sets. This is certainly a limitation, but it can be taken quite far. Yet well known facts such as Tarski’s theorem—the set of fixpoints of an isotone endofunction on a complete lattice forms a complete lattice—seem hard to formalise with it (at least without using recent experimental extensions [7]).

Firstly, leaner versions of complete lattices are introduced: Sup-lattices (and their dual Inf-lattices), in which only Sups (or Infs) are axiomatised, whereas the remaining operators, which are axiomatised in the standard Isabelle class for complete lattices, are defined explicitly. This not only reduces of proof obligations in instantiation or interpretation proofs, it also helps in constructions where only suprema are represented faithfully (e.g. using morphisms that preserve sups, but not infs, or vice versa). At the moment, Sup-lattices remain rather loosely integrated into Isabelle’s lattice hierarchy; a tighter one seems rather delicate.

Order and lattice duality is modelled, rather ad hoc, within a type class that can be added to those for orderings and lattices. Duality thus becomes a functor that reverses the order and maps Sups to Infs and vice versa, as expected. It also maps order-preserving functions to order-preserving functions, Sup-preserving to Inf-preserving ones and vice versa. This simple approach has not yet been optimised for automatic generation of dual statements (which seems hard to achieve anyway). It works quite well on simple examples.

The class-based approach to duality is contrasted by an implicit, locale-based one (which is quite standard in Isabelle), and Wenzel’s data-type-based one [11]. Wenzel’s approach generates many properties of the duality functor automatically from Isabelle’s data type package. However, duality is not involutive, and this limits the dualisation of theorems quite severely. The local-based approach dualises theorems within the context of a type class or locale highly automatically. But, unlike the present approach, it is limited to such contexts. Yet another approach to duality has been taken in HOL-Algebra [2], but it is essentially based on set theory and therefore beyond the reach of simple axiomatic type classes.

The components presented also cover fundamental concepts such as directed and filtered sets, down-closed and up-closed sets, ideals and filters, notions of sup-closure and inf-closure, sup-preservation and inf-preservation, properties of adjunctions (or Galois connections) between orderings and (complete) lattices, fusion theorems for least and greatest fixpoints, and basic properties of closure and co-closure (kernel) operations, following the Compendium (most of these concepts come as dual pairs!). As in this monograph, emphasis lies on categorical aspects, but no formal category theory is used. In addition, some simple representation theorems have been formalised, including Stone’s theorem for atomic boolean algebras (objects only). The non-atomic case seems possible, but is left for future work. Dealing with opposite maps properly, which is essential for dualities, remains an issue.

Finally, in Isabelle’s main libraries, complete distributive lattices and complete boolean algebras are currently based on a very strong distributivity law, which makes these structures *completely distributive* and is basically an Axiom of Choice. While powerset algebras satisfy this law, other appli-

cations, for instance in topology require different axiomatisations. Complete boolean algebras, in particular, are usually defined as complete lattices which are also boolean algebras. Hence only a finite distributivity law holds. Weaker distributivity laws are also essential for axiomatising complete Heyting algebras (aka frames or locales), which are relevant for point-free topology [5].

Many questions remain, in particular on tighter integrations of duality and reasoning up to isomorphism with Isabelle and beyond. In its present form, duality is often not picked up in the proofs of more complex statements. Some statements from the Compendium and Johnstone’s book had to be ignored due to the absence of carrier sets in Isabelle’s standard components for orderings and lattices. Whether Kunzar and Popescu’s new types-to-sets translation [7] provides a satisfactory solution remains to be seen.

2 Sup-Lattices and Other Simplifications

```
theory Sup-Lattice
  imports Main
begin
```

```
unbundle lattice-syntax
```

Some definitions for orderings and lattices in Isabelle could be simpler. The strict order in `ord` could be defined instead of being axiomatised. The function `mono` could have been defined on `ord` and not on `order`—even on a general (di)graph it serves as a morphism. In complete lattices, the supremum—and dually the infimum—suffices to define the other operations (in the Isabelle/HOL-definition infimum, binary supremum and infimum, bottom and top element are axiomatised). This not only increases the number of proof obligations in subclass or sublocale statements, instantiations or interpretations, it also complicates situations where suprema are presented faithfully, e.g. mapped onto suprema in some subalgebra, whereas infima in the subalgebra are different from those in the super-structure.

It would be even nicer to use a class `less-eq` which dispenses with the strict order symbol in `ord`. Then one would not have to redefine this symbol in all instantiations or interpretations. At least, it does not carry any proof obligations.

```
context ord
begin
```

`ub-set` yields the set of all upper bounds of a set; `lb-set` the set of all lower bounds.

```
definition ub-set :: 'a set ⇒ 'a set where
```

```

ub-set X = {y. ∀ x ∈ X. x ≤ y}

definition lb-set :: 'a set ⇒ 'a set where
  lb-set X = {y. ∀ x ∈ X. y ≤ x}

end

definition ord-pres :: ('a::ord ⇒ 'b::ord) ⇒ bool where
  ord-pres f = (∀ x y. x ≤ y → f x ≤ f y)

lemma ord-pres-mono:
  fixes f :: 'a::order ⇒ 'b::order
  shows mono f = ord-pres f
  by (simp add: mono-def ord-pres-def)

class preorder-lean = ord +
  assumes preorder-refl: x ≤ x
  and preorder-trans: x ≤ y ⇒ y ≤ z ⇒ x ≤ z

begin

definition le :: 'a ⇒ 'a ⇒ bool where
  le x y = (x ≤ y ∧ ¬ (x ≥ y))

end

sublocale preorder-lean ⊆ prel: preorder (≤) le
  by (unfold-locales, auto simp add: le-def preorder-refl preorder-trans)

class order-lean = preorder-lean +
  assumes order-antisym: x ≤ y ⇒ x ≥ y ⇒ x = y

sublocale order-lean ⊆ post: order (≤) le
  by (unfold-locales, simp add: order-antisym)

class Sup-lattice = order-lean + Sup +
  assumes Sups-upper: x ∈ X ⇒ x ≤ ⋒ X
  and Sups-least: (⋀ x. x ∈ X ⇒ x ≤ z) ⇒ ⋒ X ≤ z

begin

definition Infs :: 'a set ⇒ 'a where
  Infs X = ⋒ {y. ∀ x ∈ X. y ≤ x}

definition sups :: 'a ⇒ 'a ⇒ 'a where
  sups x y = ⋒ {x,y}

definition infs :: 'a ⇒ 'a ⇒ 'a where
  infs x y = Infs{x,y}

```

definition *bots* :: 'a where

$$bots = \sqcup \{\}$$

definition *tops* :: 'a where

$$tops = Infs\{\}$$

lemma *Infs-prop*: $Infs = Sup \circ lb\text{-set}$

unfolding *fun-eq-iff* **by** (*simp add: Infs-def prel.lb-set-def*)

end

class *Inf-lattice* = *order-lean* + *Inf* +

assumes *Infi-lower*: $x \in X \implies \sqcap X \leq x$

and *Infi-greatest*: $(\bigwedge x. x \in X \implies z \leq x) \implies z \leq \sqcap X$

begin

definition *Supi* :: 'a set \Rightarrow 'a where

$$Supi X = \sqcap \{y. \forall x \in X. x \leq y\}$$

definition *supi* :: 'a \Rightarrow 'a \Rightarrow 'a where

$$supi x y = Supi\{x,y\}$$

definition *infi* :: 'a \Rightarrow 'a \Rightarrow 'a where

$$infi x y = \sqcap \{x,y\}$$

definition *boti* :: 'a where

$$boti = Supi\{\}$$

definition *topi* :: 'a where

$$topi = \sqcap \{\}$$

lemma *Supi-prop*: $Supi = Inf \circ ub\text{-set}$

unfolding *fun-eq-iff* **by** (*simp add: Supi-def prel.ub-set-def*)

end

sublocale *Inf-lattice* \subseteq *ldual*: *Sup-lattice* *Inf* (\geq)

rewrites *ldual.Infs* = *Supi*

and *ldual.infs* = *supi*

and *ldual.sups* = *infi*

and *ldual.topi* = *boti*

and *ldual.bots* = *topi*

proof –

show *class.Sup-lattice* *Inf* (\geq)

by (*unfold-locales, simp-all add: Infi-lower Infi-greatest preorder-trans*)

then interpret *ldual*: *Sup-lattice* *Inf* (\geq).

show *a*: *ldual.Infs* = *Supi*

```

    unfolding fun-eq-iff by (simp add: ldual.Infs-def Supi-def)
show ldual.infs = supi
    unfolding fun-eq-iff by (simp add: a ldual.infs-def supi-def)
show ldual.sups = infi
    unfolding fun-eq-iff by (simp add: ldual.sups-def infi-def)
show ldual.tops = boti
    by (simp add: a ldual.tops-def boti-def)
show ldual.bots = topi
    by (simp add: ldual.bots-def topi-def)
qed

```

```

sublocale Sup-lattice  $\subseteq$  supclat: complete-lattice Infs Sup-class.Sup infs ( $\leq$ ) le sups
bots tops
apply unfold-locales
unfolding Infs-def infs-def sups-def bots-def tops-def
by (simp-all, auto intro: Sups-least, simp-all add: Sups-upper)

```

```

sublocale Inf-lattice  $\subseteq$  infclat: complete-lattice Inf-class.Inf Supi infi ( $\leq$ ) le supi
boti topi
by (unfold-locales, simp-all add: ldual.Sups-upper ldual.Sups-least ldual.supclat.Inf-lower
ldual.supclat.Inf-greatest)

```

end

3 Ad-Hoc Duality for Orderings and Lattices

```

theory Order-Duality
imports Sup-Lattice

```

begin

This component presents an "explicit" formalisation of order and lattice duality. It augments the data type based one used by Wenzel in his lattice components [11], and complements the "implicit" formalisation given by locales. It uses a functor *dual*, supplied within a type class, which is simply a bijection (isomorphism) between types, with the constraint that the dual of a dual object is the original object. In Wenzel's formalisation, by contrast, *dual* is a bijection, but not idempotent or involutive. In the past, Preoteasa has used a similar approach with *Isabelle* [8].

Duality is such a fundamental concept in order and lattice theory that it probably deserves to be included in the type classes for these objects, as in this section.

```

class dual =
  fixes dual :: 'a  $\Rightarrow$  'a ( $\partial$ )
  assumes inj-dual: inj  $\partial$ 
  and invol-dual [simp]:  $\partial \circ \partial = id$ 

```

This type class allows one to define a type dual. It is actually a dependent type for which dual can be instantiated.

```
typedef (overloaded) 'a dual = range (dual::'a::dual  $\Rightarrow$  'a)
  by fastforce
```

```
setup-lifting type-definition-dual
```

At the moment I have no use for this type.

```
context dual
begin
```

```
lemma invol-dual-var [simp]:  $\partial (\partial x) = x$ 
  by (simp add: pointfree-idE)
```

```
lemma surj-dual: surj  $\partial$ 
  unfolding surj-def by (metis invol-dual-var)
```

```
lemma bij-dual: bij  $\partial$ 
  by (simp add: bij-def inj-dual surj-dual)
```

```
lemma inj-dual-iff:  $(\partial x = \partial y) = (x = y)$ 
  by (meson inj-dual injD)
```

```
lemma dual-iff:  $(\partial x = y) = (x = \partial y)$ 
  by auto
```

```
lemma the-inv-dual: the-inv  $\partial = \partial$ 
  by (metis comp-apply id-def invol-dual-var inj-dual surj-dual surj-fun-eq the-inv-f-o-f-id)
```

```
end
```

In boolean algebras, duality is of course De Morgan duality and can be expressed within the language.

```
sublocale boolean-algebra  $\subseteq$  ba-dual: dual uminus
  by (unfold-locales, simp-all add: inj-def)
```

```
definition map-dual:: ('a  $\Rightarrow$  'b)  $\Rightarrow$  'a::dual  $\Rightarrow$  'b::dual ( $\partial_F$ ) where
   $\partial_F f = \partial \circ f \circ \partial$ 
```

```
lemma map-dual-func1:  $\partial_F (f \circ g) = \partial_F f \circ \partial_F g$ 
  by (metis (no-types, lifting) comp-assoc comp-id invol-dual map-dual-def)
```

```
lemma map-dual-func2 [simp]:  $\partial_F id = id$ 
  by (simp add: map-dual-def)
```

```
lemma map-dual-nat-iso:  $\partial_F f \circ \partial = \partial \circ id f$ 
  by (simp add: comp-assoc map-dual-def)
```


lemma *map-dual-invol* [*simp*]: $\partial_F \circ \partial_F = id$
unfolding *map-dual-def comp-def fun-eq-iff* **by** *simp*

Thus map-dual is naturally isomorphic to the identify functor: The function dual is a natural transformation between map-dual and the identity functor, and, because it has a two-sided inverse — itself, it is a natural isomorphism.

The generic function set-dual provides another natural transformation (see below). Before introducing it, we introduce useful notation for a widely used function.

abbreviation $\eta \equiv (\lambda x. \{x\})$

lemma *eta-inj*: *inj* η
by *simp*

definition *set-dual* = $\eta \circ \partial$

lemma *set-dual-prop*: *set-dual* (∂x) = $\{x\}$
by (*metis comp-apply dual-iff set-dual-def*)

The next four lemmas show that (functional) image and preimage are functors (on functions). This does not really belong here, but it is useful for what follows. The interaction between duality and (pre)images is needed in applications.

lemma *image-func1*: $(\cdot) (f \circ g) = (\cdot) f \circ (\cdot) g$
unfolding *fun-eq-iff* **by** (*simp add: image-comp*)

lemma *image-func2*: $(\cdot) id = id$
by *simp*

lemma *vimage-func1*: $(-\cdot) (f \circ g) = (-\cdot) g \circ (-\cdot) f$
unfolding *fun-eq-iff* **by** (*simp add: vimage-comp*)

lemma *vimage-func2*: $(-\cdot) id = id$
by *simp*

lemma *iso-image*: *mono* $((\cdot) f)$
by (*simp add: image-mono monoI*)

lemma *iso-preimage*: *mono* $((-\cdot) f)$
by (*simp add: monoI vimage-mono*)

context *dual*
begin

lemma *image-dual* [*simp*]: $(\cdot) \partial \circ (\cdot) \partial = id$
by (*metis image-func1 image-func2 invol-dual*)

lemma *vimage-dual* [simp]: $(-\cdot) \partial \circ (-\cdot) \partial = id$
by (*simp add: set.comp*)

end

The following natural transformation between the powerset functor (image) and the identity functor is well known.

lemma *power-set-func-nat-trans*: $\eta \circ id \ f = (\cdot) \ f \circ \eta$
unfolding *fun-eq-iff comp-def* **by** *simp*

As an instance, set-dual is a natural transformation with built-in type coercion.

lemma *dual-singleton*: $(\cdot) \ \partial \circ \eta = \eta \circ \partial$
by *auto*

lemma *finite-dual* [simp]: $finite \circ (\cdot) \ \partial = finite$
unfolding *fun-eq-iff comp-def* **using** *inj-dual finite-vimageI inj-vimage-image-eq*
by *fastforce*

lemma *finite-dual-var* [simp]: $finite \ (\partial \ \cdot \ X) = finite \ X$
by (*metis comp-def finite-dual*)

lemma *subset-dual*: $(X = \partial \ \cdot \ Y) = (\partial \ \cdot \ X = Y)$
by (*metis image-dual pointfree-idE*)

lemma *subset-dual1*: $(X \subseteq Y) = (\partial \ \cdot \ X \subseteq \partial \ \cdot \ Y)$
by (*simp add: inj-dual inj-image-subset-iff*)

lemma *dual-empty* [simp]: $\partial \ \cdot \ \{\} = \{\}$
by *simp*

lemma *dual-UNIV* [simp]: $\partial \ \cdot \ UNIV = UNIV$
by (*simp add: surj-dual*)

lemma *fun-dual1*: $(f = g \circ \partial) = (f \circ \partial = g)$
by (*metis comp-assoc comp-id invol-dual*)

lemma *fun-dual2*: $(f = \partial \circ g) = (\partial \circ f = g)$
by (*metis comp-assoc fun.map-id invol-dual*)

lemma *fun-dual3*: $(f = g \circ (\cdot) \ \partial) = (f \circ (\cdot) \ \partial = g)$
by (*metis comp-id image-dual o-assoc*)

lemma *fun-dual4*: $(f = (\cdot) \ \partial \circ g) = ((\cdot) \ \partial \circ f = g)$
by (*metis comp-assoc id-comp image-dual*)

lemma *fun-dual5*: $(f = \partial \circ g \circ \partial) = (\partial \circ f \circ \partial = g)$
by (*metis comp-assoc fun-dual1 fun-dual2*)

lemma *fun-dual6*: $(f = (\cdot) \partial \circ g \circ (\cdot) \partial) = ((\cdot) \partial \circ f \circ (\cdot) \partial = g)$
by (*simp add: comp-assoc fun-dual3 fun-dual4*)

lemma *fun-dual7*: $(f = \partial \circ g \circ (\cdot) \partial) = (\partial \circ f \circ (\cdot) \partial = g)$
by (*simp add: comp-assoc fun-dual2 fun-dual3*)

lemma *fun-dual8*: $(f = (\cdot) \partial \circ g \circ \partial) = ((\cdot) \partial \circ f \circ \partial = g)$
by (*simp add: comp-assoc fun-dual1 fun-dual4*)

lemma *map-dual-dual*: $(\partial_F f = g) = (\partial_F g = f)$
by (*metis map-dual-invol pointfree-idE*)

The next facts show incrementally that the dual of a complete lattice is a complete lattice.

class *ord-with-dual* = *dual* + *ord* +
assumes *ord-dual*: $x \leq y \implies \partial y \leq \partial x$

begin

lemma *dual-dual-ord*: $(\partial x \leq \partial y) = (y \leq x)$
by (*metis dual-iff ord-dual*)

end

lemma *ord-pres-dual*:
fixes $f :: 'a::ord-with-dual \Rightarrow 'b::ord-with-dual$
shows $ord-pres\ f \implies ord-pres\ (\partial_F f)$
by (*simp add: dual-dual-ord map-dual-def ord-pres-def*)

lemma *map-dual-anti*: $(f::'a::ord-with-dual \Rightarrow 'b::ord-with-dual) \leq g \implies \partial_F g \leq \partial_F f$
by (*simp add: le-fun-def map-dual-def ord-dual*)

class *preorder-with-dual* = *ord-with-dual* + *preorder*

begin

lemma *less-dual-def-var*: $(\partial y < \partial x) = (x < y)$
by (*simp add: dual-dual-ord less-le-not-le*)

end

class *order-with-dual* = *preorder-with-dual* + *order*

lemma *iso-map-dual*:
fixes $f :: 'a::order-with-dual \Rightarrow 'b::order-with-dual$
shows $mono\ f \implies mono\ (\partial_F f)$
by (*simp add: ord-pres-dual ord-pres-mono*)

```

class lattice-with-dual = lattice + dual +
  assumes sup-dual-def:  $\partial (x \sqcup y) = \partial x \sqcap \partial y$ 

begin

subclass order-with-dual
  by (unfold-locales, metis inf.absorb-iff2 sup.absorb1 sup-commute sup-dual-def)

lemma inf-dual:  $\partial (x \sqcap y) = \partial x \sqcup \partial y$ 
  by (metis invol-dual-var sup-dual-def)

lemma inf-to-sup:  $x \sqcap y = \partial (\partial x \sqcup \partial y)$ 
  using inf-dual dual-iff by fastforce

lemma sup-to-inf:  $x \sqcup y = \partial (\partial x \sqcap \partial y)$ 
  by (simp add: inf-dual)

end

class bounded-lattice-with-dual = lattice-with-dual + bounded-lattice

begin

lemma bot-dual:  $\partial \perp = \top$ 
  by (metis dual-dual-ord dual-iff le-bot top-greatest)

lemma top-dual:  $\partial \top = \perp$ 
  using bot-dual dual-iff by force

end

class boolean-algebra-with-dual = lattice-with-dual + boolean-algebra

sublocale boolean-algebra  $\subseteq$  badual: boolean-algebra-with-dual - - - - - uminus
  by unfold-locales simp-all

class Sup-lattice-with-dual = Sup-lattice + dual +
  assumes Sups-dual-def:  $\partial \circ \text{Sup} = \text{Infs} \circ (\cdot) \partial$ 

class Inf-lattice-with-dual = Inf-lattice + dual +
  assumes Sups-dual-def:  $\partial \circ \text{Supi} = \text{Inf} \circ (\cdot) \partial$ 

class complete-lattice-with-dual = complete-lattice + dual +
  assumes Sups-dual-def:  $\partial \circ \text{Sup} = \text{Inf} \circ (\cdot) \partial$ 

sublocale Sup-lattice-with-dual  $\subseteq$  sclatd: complete-lattice-with-dual Infs Sup infs
  ( $\leq$ ) le sups bots tops  $\partial$ 
  by (unfold-locales, simp add: Sups-dual-def)

```

sublocale *Inf-lattice-with-dual* \subseteq *iclatd: complete-lattice-with-dual* *Inf Supi infi*
 (\leq) *le supi boti topi* ∂
by (*unfold-locales, simp add: Sups-dual-def*)

context *complete-lattice-with-dual*
begin

lemma *Inf-dual*: $\partial \circ \text{Inf} = \text{Sup} \circ (\cdot) \partial$
by (*metis comp-assoc comp-id fun.map-id Sups-dual-def image-dual invol-dual*)

lemma *Inf-dual-var*: $\partial (\prod X) = \bigsqcup (\partial \cdot X)$
using *comp-eq-dest Inf-dual* **by** *fastforce*

lemma *Inf-to-Sup*: $\text{Inf} = \partial \circ \text{Sup} \circ (\cdot) \partial$
by (*auto simp add: Sups-dual-def image-comp*)

lemma *Inf-to-Sup-var*: $\prod X = \partial (\bigsqcup (\partial \cdot X))$
using *Inf-dual-var dual-iff* **by** *fastforce*

lemma *Sup-to-Inf*: $\text{Sup} = \partial \circ \text{Inf} \circ (\cdot) \partial$
by (*auto simp add: Inf-dual image-comp*)

lemma *Sup-to-Inf-var*: $\bigsqcup X = \partial (\prod (\partial \cdot X))$
using *Sup-to-Inf* **by** *force*

lemma *Sup-dual-def-var*: $\partial (\bigsqcup X) = \prod (\partial \cdot X)$
using *comp-eq-dest Sups-dual-def* **by** *fastforce*

lemma *bot-dual-def*: $\partial \top = \perp$
by (*smt (verit) Inf-UNIV Sup-UNIV Sups-dual-def surj-dual o-eq-dest*)

lemma *top-dual-def*: $\partial \perp = \top$
using *bot-dual-def dual-iff* **by** *blast*

lemma *inf-dual2*: $\partial (x \sqcap y) = \partial x \sqcup \partial y$
by (*smt (verit) comp-eq-elim Inf-dual Inf-empty Inf-insert SUP-insert inf-top.right-neutral*)

lemma *sup-dual*: $\partial (x \sqcup y) = \partial x \sqcap \partial y$
by (*metis inf-dual2 dual-iff*)

subclass *lattice-with-dual*
by (*unfold-locales, auto simp: inf-dual sup-dual*)

subclass *bounded-lattice-with-dual..*

end

end

4 Properties of Orderings and Lattices

theory *Order-Lattice-Props*
imports *Order-Duality*

begin

4.1 Basic Definitions for Orderings and Lattices

The first definition is for order morphisms — isotone (order-preserving, monotone) functions. An order isomorphism is an order-preserving bijection. This should be defined in the class *ord*, but *mono* requires order.

definition *ord-homset* :: ('a::order \Rightarrow 'b::order) set **where**
ord-homset = {f::'a::order \Rightarrow 'b::order. *mono* f}

definition *ord-embed* :: ('a::order \Rightarrow 'b::order) \Rightarrow bool **where**
ord-embed f = ($\forall x y. f x \leq f y \longleftrightarrow x \leq y$)

definition *ord-iso* :: ('a::order \Rightarrow 'b::order) \Rightarrow bool **where**
ord-iso = *bij* \sqcap *mono* \sqcap (*mono* \circ *the-inv*)

lemma *ord-embed-alt*: *ord-embed* f = (*mono* f \wedge ($\forall x y. f x \leq f y \longrightarrow x \leq y$))
using *mono-def ord-embed-def* **by** *auto*

lemma *ord-embed-homset*: *ord-embed* f \Longrightarrow f \in *ord-homset*
by (*simp add: mono-def ord-embed-def ord-homset-def*)

lemma *ord-embed-inj*: *ord-embed* f \Longrightarrow *inj* f
unfolding *ord-embed-def inj-def* **by** (*simp add: eq-iff*)

lemma *ord-iso-ord-embed*: *ord-iso* f \Longrightarrow *ord-embed* f
unfolding *ord-iso-def ord-embed-def bij-def inj-def mono-def*
by (*clarsimp, metis inj-def the-inv-f-f*)

lemma *ord-iso-alt*: *ord-iso* f = (*ord-embed* f \wedge *surj* f)
unfolding *ord-iso-def ord-embed-def surj-def bij-def inj-def mono-def*
apply *safe*
by *simp-all (metis eq-iff inj-def the-inv-f-f)+*

lemma *ord-iso-the-inv*: *ord-iso* f \Longrightarrow *mono* (*the-inv* f)
by (*simp add: ord-iso-def*)

lemma *ord-iso-inv1*: *ord-iso* f \Longrightarrow (*the-inv* f) \circ f = *id*
using *ord-embed-inj ord-iso-ord-embed the-inv-into-f-f* **by** *fastforce*

lemma *ord-iso-inv2*: *ord-iso* f \Longrightarrow f \circ (*the-inv* f) = *id*
using *f-the-inv-into-f ord-embed-inj ord-iso-alt* **by** *fastforce*

typedef (overloaded) ('a,'b) ord-homset = ord-homset::('a::order \Rightarrow 'b::order)
set

by (force simp: ord-homset-def mono-def)

setup-lifting type-definition-ord-homset

The next definition is for the set of fixpoints of a given function. It is important in the context of orders, for instance for proving Tarski's fixpoint theorem, but does not really belong here.

definition Fix :: ('a \Rightarrow 'a) \Rightarrow 'a set **where**

Fix f = {x. f x = x}

lemma retraction-prop: f \circ f = f \implies f x = x \longleftrightarrow x \in range f

by (metis comp-apply f-inv-into-f rangeI)

lemma retraction-prop-fix: f \circ f = f \implies range f = Fix f

unfolding Fix-def **using** retraction-prop **by** fastforce

lemma Fix-map-dual: Fix \circ ∂_F = (\wedge) ∂ \circ Fix

unfolding Fix-def map-dual-def comp-def fun-eq-iff

by (smt (verit) Collect-cong invol-dual pointfree-idE setcompr-eq-image)

lemma Fix-map-dual-var: Fix (∂_F f) = ∂ ' (Fix f)

by (metis Fix-map-dual o-def)

lemma gfp-dual: (∂ ::'a::complete-lattice-with-dual \Rightarrow 'a) \circ gfp = lfp \circ ∂_F

proof –

{**fix** f:: 'a \Rightarrow 'a

have ∂ (gfp f) = ∂ (\bigsqcup {u. u \leq f u})

by (simp add: gfp-def)

also have ... = \bigsqcap (∂ ' {u. u \leq f u})

by (simp add: Sup-dual-def-var)

also have ... = \bigsqcap { ∂ u | u. u \leq f u}

by (simp add: setcompr-eq-image)

also have ... = \bigsqcap {u | u. (∂_F f) u \leq u}

by (metis (no-types, opaque-lifting) dual-dual-ord dual-iff map-dual-def o-def)

finally have ∂ (gfp f) = lfp (∂_F f)

by (metis lfp-def)}

thus ?thesis

by auto

qed

lemma gfp-dual-var:

fixes f :: 'a::complete-lattice-with-dual \Rightarrow 'a

shows ∂ (gfp f) = lfp (∂_F f)

using comp-eq-elim gfp-dual **by** blast

lemma gfp-to-lfp: gfp = (∂ ::'a::complete-lattice-with-dual \Rightarrow 'a) \circ lfp \circ ∂_F

by (simp add: comp-assoc fun-dual2 gfp-dual)

lemma *gfp-to-lfp-var*:
fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$
shows $\text{gfp } f = \partial (\text{lfp } (\partial_F f))$
by (*metis* *gfp-dual-var invol-dual-var*)

lemma *lfp-dual*: $(\partial::'a::\text{complete-lattice-with-dual} \Rightarrow 'a) \circ \text{lfp} = \text{gfp} \circ \partial_F$
by (*simp add: comp-assoc gfp-to-lfp map-dual-invol*)

lemma *lfp-dual-var*:
fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$
shows $\partial (\text{lfp } f) = \text{gfp } (\text{map-dual } f)$
using *comp-eq-dest-lhs lfp-dual* **by** *fastforce*

lemma *lfp-to-gfp*: $\text{lfp} = (\partial::'a::\text{complete-lattice-with-dual} \Rightarrow 'a) \circ \text{gfp} \circ \partial_F$
by (*simp add: comp-assoc gfp-dual map-dual-invol*)

lemma *lfp-to-gfp-var*:
fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$
shows $\text{lfp } f = \partial (\text{gfp } (\partial_F f))$
by (*metis* *invol-dual-var lfp-dual-var*)

lemma *lfp-in-Fix*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'a$
shows $\text{mono } f \Longrightarrow \text{lfp } f \in \text{Fix } f$
by (*metis* (*mono-tags, lifting*) *Fix-def lfp-unfold mem-Collect-eq*)

lemma *gfp-in-Fix*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'a$
shows $\text{mono } f \Longrightarrow \text{gfp } f \in \text{Fix } f$
by (*metis* (*mono-tags, lifting*) *Fix-def gfp-unfold mem-Collect-eq*)

lemma *nonempty-Fix*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'a$
shows $\text{mono } f \Longrightarrow \text{Fix } f \neq \{\}$
using *lfp-in-Fix* **by** *fastforce*

Next the minimal and maximal elements of an ordering are defined.

context *ord*
begin

definition *min-set* :: $'a \text{ set} \Rightarrow 'a \text{ set}$ **where**
 $\text{min-set } X = \{y \in X. \forall x \in X. x \leq y \longrightarrow x = y\}$

definition *max-set* :: $'a \text{ set} \Rightarrow 'a \text{ set}$ **where**
 $\text{max-set } X = \{x \in X. \forall y \in X. x \leq y \longrightarrow x = y\}$

end

context *ord-with-dual*
begin

lemma *min-max-set-dual*: $(\cdot) \partial \circ \text{min-set} = \text{max-set} \circ (\cdot) \partial$
unfolding *max-set-def min-set-def fun-eq-iff comp-def*
apply *safe*
using *dual-dual-ord inj-dual-iff* **by** *auto*

lemma *min-max-set-dual-var*: $\partial \cdot (\text{min-set } X) = \text{max-set } (\partial \cdot X)$
using *comp-eq-dest min-max-set-dual* **by** *fastforce*

lemma *max-min-set-dual*: $(\cdot) \partial \circ \text{max-set} = \text{min-set} \circ (\cdot) \partial$
by (*metis (no-types, opaque-lifting) comp-id fun.map-comp id-comp image-dual min-max-set-dual*)

lemma *min-to-max-set*: $\text{min-set} = (\cdot) \partial \circ \text{max-set} \circ (\cdot) \partial$
by (*metis comp-id image-dual max-min-set-dual o-assoc*)

lemma *max-min-set-dual-var*: $\partial \cdot (\text{max-set } X) = \text{min-set } (\partial \cdot X)$
using *comp-eq-dest max-min-set-dual* **by** *fastforce*

lemma *min-to-max-set-var*: $\text{min-set } X = \partial \cdot (\text{max-set } (\partial \cdot X))$
by (*simp add: max-min-set-dual-var pointfree-idE*)

end

Next, directed and filtered sets, upsets, downsets, filters and ideals in posets are defined.

context *ord*
begin

definition *directed* :: *'a set* \Rightarrow *bool* **where**
 $\text{directed } X = (\forall Y. \text{finite } Y \wedge Y \subseteq X \longrightarrow (\exists x \in X. \forall y \in Y. y \leq x))$

definition *filtered* :: *'a set* \Rightarrow *bool* **where**
 $\text{filtered } X = (\forall Y. \text{finite } Y \wedge Y \subseteq X \longrightarrow (\exists x \in X. \forall y \in Y. x \leq y))$

definition *downset-set* :: *'a set* \Rightarrow *'a set* (\Downarrow) **where**
 $\Downarrow X = \{y. \exists x \in X. y \leq x\}$

definition *upset-set* :: *'a set* \Rightarrow *'a set* (\Uparrow) **where**
 $\Uparrow X = \{y. \exists x \in X. x \leq y\}$

definition *downset* :: *'a* \Rightarrow *'a set* (\downarrow) **where**
 $\downarrow = \Downarrow \circ \eta$

definition *upset* :: *'a* \Rightarrow *'a set* (\uparrow) **where**
 $\uparrow = \Uparrow \circ \eta$

definition *downsets* :: 'a set set **where**
downsets = *Fix* \Downarrow

definition *upsets* :: 'a set set **where**
upsets = *Fix* \Uparrow

definition *downclosed-set* $X = (X \in \text{downsets})$

definition *upclosed-set* $X = (X \in \text{upsets})$

definition *ideals* :: 'a set set **where**
ideals = $\{X. X \neq \{\} \wedge \text{downclosed-set } X \wedge \text{directed } X\}$

definition *filters* :: 'a set set **where**
filters = $\{X. X \neq \{\} \wedge \text{upclosed-set } X \wedge \text{filtered } X\}$

abbreviation *idealp* $X \equiv X \in \text{ideals}$

abbreviation *filterp* $X \equiv X \in \text{filters}$

end

These notions are pair-wise dual.

Filtered and directed sets are dual.

context *ord-with-dual*
begin

lemma *filtered-directed-dual*: *filtered* $\circ (\cdot) \partial = \text{directed}$
unfolding *filtered-def directed-def fun-eq-iff comp-def*
apply *clarsimp*
apply *safe*
apply (*meson finite-imageI imageI image-mono dual-dual-ord*)
by (*smt (verit, ccfv-threshold) finite-subset-image imageE ord-dual*)

lemma *directed-filtered-dual*: *directed* $\circ (\cdot) \partial = \text{filtered}$
using *filtered-directed-dual* **by** (*metis comp-id image-dual o-assoc*)

lemma *filtered-to-directed*: *filtered* $X = \text{directed } (\partial \text{ ` } X)$
by (*metis comp-apply directed-filtered-dual*)

Upsets and downsets are dual.

lemma *downset-set-upset-set-dual*: $(\cdot) \partial \circ \Downarrow = \Uparrow \circ (\cdot) \partial$
unfolding *downset-set-def upset-set-def fun-eq-iff comp-def*
apply *safe*
apply (*meson image-eqI ord-dual*)
by (*clarsimp, metis (mono-tags, lifting) dual-iff image-iff mem-Collect-eq ord-dual*)

lemma *upset-set-downset-set-dual*: $(\cdot) \partial \circ \Uparrow = \Downarrow \circ (\cdot) \partial$

using *downset-set-upset-set-dual* **by** (*metis (no-types, opaque-lifting) comp-id id-comp image-dual o-assoc*)

lemma *upset-set-to-downset-set*: $\uparrow = (\cdot) \partial \circ \downarrow \circ (\cdot) \partial$
by (*simp add: comp-assoc downset-set-upset-set-dual*)

lemma *upset-set-to-downset-set2*: $\uparrow X = \partial \cdot (\downarrow (\partial \cdot X))$
by (*simp add: upset-set-to-downset-set*)

lemma *downset-upset-dual*: $(\cdot) \partial \circ \downarrow = \uparrow \circ \partial$
using *downset-def upset-def upset-set-to-downset-set* **by** *fastforce*

lemma *upset-to-downset*: $(\cdot) \partial \circ \uparrow = \downarrow \circ \partial$
by (*metis comp-assoc id-apply ord.downset-def ord.upset-def power-set-func-nat-trans upset-set-downset-set-dual*)

lemma *upset-to-downset2*: $\uparrow = (\cdot) \partial \circ \downarrow \circ \partial$
by (*simp add: comp-assoc downset-upset-dual*)

lemma *upset-to-downset3*: $\uparrow x = \partial \cdot (\downarrow (\partial x))$
by (*simp add: upset-to-downset2*)

lemma *downsets-upsets-dual*: $(X \in \text{downsets}) = (\partial \cdot X \in \text{upsets})$
unfolding *downsets-def upsets-def Fix-def*
by (*smt (verit) comp-eq-dest downset-set-upset-set-dual image-inv-f-f inj-dual mem-Collect-eq*)

lemma *downset-setp-upset-setp-dual*: $\text{upclosed-set} \circ (\cdot) \partial = \text{downclosed-set}$
unfolding *downclosed-set-def upclosed-set-def* **using** *downsets-upsets-dual* **by** *fastforce*

lemma *upsets-to-downsets*: $(X \in \text{upsets}) = (\partial \cdot X \in \text{downsets})$
by (*simp add: downsets-upsets-dual image-comp*)

lemma *upset-setp-downset-setp-dual*: $\text{downclosed-set} \circ (\cdot) \partial = \text{upclosed-set}$
by (*metis comp-id downset-setp-upset-setp-dual image-dual o-assoc*)

Filters and ideals are dual.

lemma *ideals-filters-dual*: $(X \in \text{ideals}) = ((\partial \cdot X) \in \text{filters})$
by (*smt (verit) comp-eq-dest-lhs directed-filtered-dual image-inv-f-f image-is-empty inv-unique-comp filters-def ideals-def inj-dual invol-dual mem-Collect-eq upset-setp-downset-setp-dual*)

lemma *idealp-filterp-dual*: $\text{idealp} = \text{filterp} \circ (\cdot) \partial$
unfolding *fun-eq-iff* **by** (*simp add: ideals-filters-dual*)

lemma *filters-to-ideals*: $(X \in \text{filters}) = ((\partial \cdot X) \in \text{ideals})$
by (*simp add: ideals-filters-dual image-comp*)

lemma *filterp-idealp-dual*: $\text{filterp} = \text{idealp} \circ (\cdot) \partial$

```

unfolding fun-eq-iff by (simp add: filters-to-ideals)

end

## 4.2 Properties of Orderings

context ord
begin

lemma directed-nonempty: directed  $X \implies X \neq \{\}$ 
  unfolding directed-def by fastforce

lemma directed-ub: directed  $X \implies (\forall x \in X. \forall y \in X. \exists z \in X. x \leq z \wedge y \leq z)$ 
  by (meson empty-subsetI directed-def finite.emptyI finite-insert insert-subset order-refl)

lemma downset-set-prop:  $\Downarrow = \text{Union} \circ (\cdot) \downarrow$ 
  unfolding downset-set-def downset-def fun-eq-iff by fastforce

lemma downset-set-prop-var:  $\Downarrow X = (\bigcup x \in X. \downarrow x)$ 
  by (simp add: downset-set-prop)

lemma downset-prop:  $\downarrow x = \{y. y \leq x\}$ 
  unfolding downset-def downset-set-def fun-eq-iff by fastforce

lemma downset-prop2:  $y \leq x \implies y \in \downarrow x$ 
  by (simp add: downset-prop)

lemma ideals-downsets:  $X \in \text{ideals} \implies X \in \text{downsets}$ 
  by (simp add: downclosed-set-def ideals-def)

lemma ideals-directed:  $X \in \text{ideals} \implies \text{directed } X$ 
  by (simp add: ideals-def)

end

context preorder
begin

lemma directed-prop:  $X \neq \{\} \implies (\forall x \in X. \forall y \in X. \exists z \in X. x \leq z \wedge y \leq z) \implies \text{directed } X$ 
proof –
  assume h1:  $X \neq \{\}$ 
  and h2:  $\forall x \in X. \forall y \in X. \exists z \in X. x \leq z \wedge y \leq z$ 
  {fix Y
  have finite  $Y \implies Y \subseteq X \implies (\exists x \in X. \forall y \in Y. y \leq x)$ 
  proof (induct rule: finite-induct)
    case empty
    then show ?case

```

```

    using h1 by blast
  next
  case (insert x F)
  then show ?case
    by (metis h2 insert-iff insert-subset order-trans)
  qed}
  thus ?thesis
    by (simp add: directed-def)
qed

lemma directed-alt: directed X = (X ≠ {} ∧ (∀ x ∈ X. ∀ y ∈ X. ∃ z ∈ X. x ≤ z
  ∧ y ≤ z))
  by (metis directed-prop directed-nonempty directed-ub)

lemma downset-set-prop-var2: x ∈ ↓X ⇒ y ≤ x ⇒ y ∈ ↓X
  unfolding downset-set-def using order-trans by blast

lemma downclosed-set-iff: downclosed-set X = (∀ x ∈ X. ∀ y. y ≤ x → y ∈ X)
  unfolding downclosed-set-def downsets-def Fix-def downset-set-def by auto

lemma downclosed-downset-set: downclosed-set (↓X)
  by (simp add: downclosed-set-iff downset-set-prop-var2 downset-def)

lemma downclosed-downset: downclosed-set (↓x)
  by (simp add: downclosed-downset-set downset-def)

lemma downset-set-ext: id ≤ ↓
  unfolding le-fun-def id-def downset-set-def by auto

lemma downset-set-iso: mono ↓
  unfolding mono-def downset-set-def by blast

lemma downset-set-idem [simp]: ↓ ∘ ↓ = ↓
  unfolding fun-eq-iff downset-set-def using order-trans by auto

lemma downset-faithful: ↓x ⊆ ↓y ⇒ x ≤ y
  by (simp add: downset-prop subset-eq)

lemma downset-iso-iff: (↓x ⊆ ↓y) = (x ≤ y)
  using atMost-iff downset-prop order-trans by blast

```

The following proof uses the Axiom of Choice.

```

lemma downset-directed-downset-var [simp]: directed (↓X) = directed X
proof
  assume h1: directed X
  {fix Y
  assume h2: finite Y and h3: Y ⊆ ↓X
  hence ∀ y. ∃ x. y ∈ Y → x ∈ X ∧ y ≤ x
    by (force simp: downset-set-def)
  }

```

hence $\exists f. \forall y. y \in Y \longrightarrow f y \in X \wedge y \leq f y$
by (*rule choice*)
hence $\exists f. \text{finite } (f \text{ ` } Y) \wedge f \text{ ` } Y \subseteq X \wedge (\forall y \in Y. y \leq f y)$
by (*metis finite-imageI h2 image-subsetI*)
hence $\exists Z. \text{finite } Z \wedge Z \subseteq X \wedge (\forall y \in Y. \exists z \in Z. y \leq z)$
by *fastforce*
hence $\exists Z. \text{finite } Z \wedge Z \subseteq X \wedge (\forall y \in Y. \exists z \in Z. y \leq z) \wedge (\exists x \in X. \forall z \in Z. z \leq x)$
by (*metis directed-def h1*)
hence $\exists x \in X. \forall y \in Y. y \leq x$
by (*meson order-trans*)
thus *directed* $(\Downarrow X)$
unfolding *directed-def downset-set-def* **by** *fastforce*
next
assume *directed* $(\Downarrow X)$
thus *directed* X
unfolding *directed-def downset-set-def*
apply *clarsimp*
by (*smt (verit) Ball-Collect order-refl order-trans subsetCE*)
qed

lemma *downset-directed-downset* [*simp*]: $\text{directed} \circ \Downarrow = \text{directed}$
unfolding *fun-eq-iff* **by** *simp*

lemma *directed-downset-ideals*: $\text{directed } (\Downarrow X) = (\Downarrow X \in \text{ideals})$
by (*metis (mono-tags, lifting) CollectI Fix-def directed-alt downset-set-idem down-closed-set-def downsets-def ideals-def o-def ord.ideals-directed*)

lemma *downclosed-Fix*: $\text{downclosed-set } X = (\Downarrow X = X)$
by (*metis (mono-tags, lifting) CollectD Fix-def downclosed-downset-set down-closed-set-def downsets-def*)

end

lemma *downset-iso*: $\text{mono } (\Downarrow :: 'a::\text{order} \Rightarrow 'a \text{ set})$
by (*simp add: downset-iso-iff mono-def*)

lemma *mono-downclosed*:
fixes $f :: 'a::\text{order} \Rightarrow 'b::\text{order}$
assumes *mono* f
shows $\forall Y. \text{downclosed-set } Y \longrightarrow \text{downclosed-set } (f \text{ ` } Y)$
by (*simp add: asms downclosed-set-iff monoD*)

lemma
fixes $f :: 'a::\text{order} \Rightarrow 'b::\text{order}$
assumes *mono* f
shows $\forall Y. \text{downclosed-set } X \longrightarrow \text{downclosed-set } (f \text{ ` } X)$
oops

```

lemma downclosed-mono:
  fixes  $f :: 'a::order \Rightarrow 'b::order$ 
  assumes  $\forall Y. \text{downclosed-set } Y \longrightarrow \text{downclosed-set } (f -' Y)$ 
  shows mono f
proof -
  {fix  $x y :: 'a::order$ 
  assume  $h: x \leq y$ 
  have downclosed-set  $(\downarrow (f y))$ 
    unfolding downclosed-set-def downsets-def Fix-def downset-set-def downset-def
  by auto
  hence downclosed-set  $(f -' (\downarrow (f y)))$ 
    by (simp add: assms)
  hence downclosed-set  $\{z. f z \leq f y\}$ 
    unfolding vimage-def downset-def downset-set-def by auto
  hence  $\forall z w. (f z \leq f y \wedge w \leq z) \longrightarrow f w \leq f y$ 
    unfolding downclosed-set-def downclosed-set-def downsets-def Fix-def downset-set-def
  by force
  hence  $f x \leq f y$ 
    using  $h$  by blast}
  thus ?thesis..
qed

```

```

lemma mono-downclosed-iff:  $\text{mono } f = (\forall Y. \text{downclosed-set } Y \longrightarrow \text{downclosed-set } (f -' Y))$ 
using mono-downclosed downclosed-mono by auto

```

```

context order
begin

```

```

lemma downset-inj:  $\text{inj } \downarrow$ 
by (metis injI downset-iso-iff order.eq-iff)

```

```

lemma  $(X \subseteq Y) = (\downarrow X \subseteq \downarrow Y)$ 
oops

```

```

end

```

```

context lattice
begin

```

```

lemma lat-ideals:  $X \in \text{ideals} = (X \neq \{\} \wedge X \in \text{downsets} \wedge (\forall x \in X. \forall y \in X. x \sqcup y \in X))$ 
unfolding ideals-def directed-alt downsets-def Fix-def downset-set-def downclosed-set-def
using local.sup.bounded-iff local.sup-ge2 by blast

```

```

end

```

```

context bounded-lattice

```

begin

lemma *bot-ideal*: $X \in \text{ideals} \implies \perp \in X$

unfolding *ideals-def downclosed-set-def downsets-def Fix-def downset-set-def* **by** *fastforce*

end

context *complete-lattice*

begin

lemma *Sup-downset-id* [*simp*]: $\text{Sup} \circ \downarrow = \text{id}$

using *Sup-atMost atMost-def downset-prop* **by** *fastforce*

lemma *downset-Sup-id*: $\text{id} \leq \downarrow \circ \text{Sup}$

by (*simp add: Sup-upper downset-prop le-funI subsetI*)

lemma *Inf-Sup-var*: $\bigsqcup (\bigcap x \in X. \downarrow x) = \bigsqcap X$

unfolding *downset-prop* **by** (*simp add: Collect-ball-eq Inf-eq-Sup*)

lemma *Inf-pres-downset-var*: $(\bigcap x \in X. \downarrow x) = \downarrow(\bigsqcap X)$

unfolding *downset-prop* **by** (*safe, simp-all add: le-Inf-iff*)

end

4.3 Dual Properties of Orderings

context *ord-with-dual*

begin

lemma *filtered-nonempty*: $\text{filtered } X \implies X \neq \{\}$

using *filtered-to-directed ord.directed-nonempty* **by** *auto*

lemma *filtered-lb*: $\text{filtered } X \implies (\forall x \in X. \forall y \in X. \exists z \in X. z \leq x \wedge z \leq y)$

using *filtered-to-directed directed-ub dual-dual-ord* **by** *fastforce*

lemma *upset-set-prop-var*: $\uparrow X = (\bigcup x \in X. \uparrow x)$

by (*simp add: image-Union downset-set-prop-var upset-set-to-downset-set2 upset-to-downset2*)

lemma *upset-set-prop*: $\uparrow = \text{Union} \circ (\cdot) \uparrow$

unfolding *fun-eq-iff* **by** (*simp add: upset-set-prop-var*)

lemma *upset-prop*: $\uparrow x = \{y. x \leq y\}$

unfolding *upset-to-downset3 downset-prop image-def* **using** *dual-dual-ord* **by** *fastforce*

lemma *upset-prop2*: $x \leq y \implies y \in \uparrow x$

by (*simp add: upset-prop*)

lemma *filters-upsets*: $X \in \text{filters} \implies X \in \text{upsets}$
by (*simp add: upclosed-set-def filters-def*)

lemma *filters-filtered*: $X \in \text{filters} \implies \text{filtered } X$
by (*simp add: filters-def*)

end

context *preorder-with-dual*
begin

lemma *filtered-prop*: $X \neq \{\}$ $\implies (\forall x \in X. \forall y \in X. \exists z \in X. z \leq x \wedge z \leq y) \implies$
filtered X
unfolding *filtered-to-directed*
by (*rule directed-prop, blast, metis (full-types) image-iff ord-dual*)

lemma *filtered-alt*: $\text{filtered } X = (X \neq \{\}) \wedge (\forall x \in X. \forall y \in X. \exists z \in X. z \leq x \wedge z \leq y)$
by (*metis image-empty directed-alt filtered-to-directed filtered-lb filtered-prop*)

lemma *up-set-prop-var2*: $x \in \uparrow X \implies x \leq y \implies y \in \uparrow X$
using *downset-set-prop-var2 dual-iff ord-dual upset-set-to-downset-set2* **by** *fast-force*

lemma *upclosed-set-iff*: $\text{upclosed-set } X = (\forall x \in X. \forall y. x \leq y \longrightarrow y \in X)$
unfolding *upclosed-set-def upsets-def Fix-def upset-set-def* **by** *auto*

lemma *upclosed-upset-set*: $\text{upclosed-set } (\uparrow X)$
using *up-set-prop-var2 upclosed-set-iff* **by** *blast*

lemma *upclosed-upset*: $\text{upclosed-set } (\uparrow x)$
by (*simp add: upset-def upclosed-upset-set*)

lemma *upset-set-ext*: $\text{id} \leq \uparrow$
by (*smt (verit) comp-def comp-id image-mono le-fun-def downset-set-ext image-dual upset-set-to-downset-set2*)

lemma *upset-set-anti*: $\text{mono } \uparrow$
by (*metis image-mono downset-set-iso upset-set-to-downset-set2 mono-def*)

lemma *up-set-idem* [*simp*]: $\uparrow \circ \uparrow = \uparrow$
by (*metis comp-assoc downset-set-idem upset-set-downset-set-dual upset-set-to-downset-set*)

lemma *upset-faithful*: $\uparrow x \subseteq \uparrow y \implies y \leq x$
by (*metis inj-image-subset-iff downset-faithful dual-dual-ord inj-dual upset-to-downset3*)

lemma *upset-anti-iff*: $(\uparrow y \subseteq \uparrow x) = (x \leq y)$
by (*metis downset-iso-iff ord-dual upset-to-downset3 subset-image-iff upset-faithful*)

lemma *upset-filtered-upset* [*simp*]: $\text{filtered} \circ \uparrow = \text{filtered}$
by (*metis comp-assoc directed-filtered-dual downset-directed-downset upset-set-downset-set-dual*)

lemma *filtered-upset-filters*: $\text{filtered} (\uparrow X) = (\uparrow X \in \text{filters})$
by (*metis comp-apply directed-downset-ideals filtered-to-directed filterp-idealp-dual upset-set-downset-set-dual*)

lemma *upclosed-Fix*: $\text{upclosed-set } X = (\uparrow X = X)$
by (*simp add: Fix-def upclosed-set-def upsets-def*)

end

lemma *upset-anti*: $\text{antimono} (\uparrow :: 'a::\text{order-with-dual} \Rightarrow 'a \text{ set})$
by (*simp add: antimono-def upset-anti-iff*)

lemma *mono-upclosed*:
fixes $f :: 'a::\text{order-with-dual} \Rightarrow 'b::\text{order-with-dual}$
assumes *mono f*
shows $\forall Y. \text{upclosed-set } Y \longrightarrow \text{upclosed-set } (f -' Y)$
by (*simp add: assms monoD upclosed-set-iff*)

lemma *mono-upclosed*:
fixes $f :: 'a::\text{order-with-dual} \Rightarrow 'b::\text{order-with-dual}$
assumes *mono f*
shows $\forall Y. \text{upclosed-set } X \longrightarrow \text{upclosed-set } (f ' X)$
oops

lemma *upclosed-mono*:
fixes $f :: 'a::\text{order-with-dual} \Rightarrow 'b::\text{order-with-dual}$
assumes $\forall Y. \text{upclosed-set } Y \longrightarrow \text{upclosed-set } (f -' Y)$
shows *mono f*
by (*metis (mono-tags, lifting) assms dual-order.refl mem-Collect-eq monoI order.trans upclosed-set-iff vimageE vimageI2*)

lemma *mono-upclosed-iff*:
fixes $f :: 'a::\text{order-with-dual} \Rightarrow 'b::\text{order-with-dual}$
shows $\text{mono } f = (\forall Y. \text{upclosed-set } Y \longrightarrow \text{upclosed-set } (f -' Y))$
using *mono-upclosed upclosed-mono* **by** *auto*

context *order-with-dual*
begin

lemma *upset-inj*: $\text{inj } \uparrow$
by (*metis inj-compose inj-on-imageI2 downset-inj inj-dual upset-to-downset*)

lemma $(X \subseteq Y) = (\uparrow Y \subseteq \uparrow X)$
oops

end

context *lattice-with-dual*
begin

lemma *lat-filters*: $X \in \text{filters} = (X \neq \{\}) \wedge X \in \text{upsets} \wedge (\forall x \in X. \forall y \in X. x \sqcap y \in X)$

unfolding *filters-to-ideals upsets-to-downsets inf-to-sup lat-ideals*

by (*smt (verit) image-iff image-inv-f-f image-is-empty inj-image-mem-iff inv-unique-comp inj-dual invol-dual*)

end

context *bounded-lattice-with-dual*
begin

lemma *top-filter*: $X \in \text{filters} \implies \top \in X$

using *bot-ideal inj-image-mem-iff inj-dual filters-to-ideals top-dual* **by** *fastforce*

end

context *complete-lattice-with-dual*
begin

lemma *Inf-upset-id [simp]*: $\text{Inf} \circ \uparrow = \text{id}$

by (*metis comp-assoc comp-id Sup-downset-id Sups-dual-def downset-upset-dual invol-dual*)

lemma *upset-Inf-id*: $\text{id} \leq \uparrow \circ \text{Inf}$

by (*simp add: Inf-lower le-funI subsetI upset-prop*)

lemma *Sup-Inf-var*: $\sqcap (\bigcap x \in X. \uparrow x) = \sqcup X$

unfolding *upset-prop* **by** (*simp add: Collect-ball-eq Sup-eq-Inf*)

lemma *Sup-dual-upset-var*: $(\bigcap x \in X. \uparrow x) = \uparrow(\sqcup X)$

unfolding *upset-prop* **by** (*safe, simp-all add: Sup-le-iff*)

end

4.4 Properties of Complete Lattices

definition *Inf-closed-set* $X = (\forall Y \subseteq X. \sqcap Y \in X)$

definition *Sup-closed-set* $X = (\forall Y \subseteq X. \sqcup Y \in X)$

definition *inf-closed-set* $X = (\forall x \in X. \forall y \in X. x \sqcap y \in X)$

definition *sup-closed-set* $X = (\forall x \in X. \forall y \in X. x \sqcup y \in X)$

The following facts about complete lattices add to those in the Isabelle

libraries.

context *complete-lattice*
begin

The translation between sup and Sup could be improved. The sup-theorems should be direct consequences of Sup-ones. In addition, duality between sup and inf is currently not exploited.

lemma *sup-Sup*: $x \sqcup y = \sqcup \{x, y\}$
by *simp*

lemma *inf-Inf*: $x \sqcap y = \sqcap \{x, y\}$
by *simp*

The next two lemmas are about Sups and Infs of indexed families. These are interesting for iterations and fixpoints.

lemma *fSup-unfold*: $(f :: nat \Rightarrow 'a) \ 0 \sqcup (\sqcup n. f (Suc n)) = (\sqcup n. f n)$
apply (*intro order.antisym sup-least*)
apply (*rule Sup-upper, force*)
apply (*rule Sup-mono, force*)
apply (*safe intro!: Sup-least*)
by (*case-tac n, simp-all add: Sup-upper le-supI2*)

lemma *fInf-unfold*: $(f :: nat \Rightarrow 'a) \ 0 \sqcap (\sqcap n. f (Suc n)) = (\sqcap n. f n)$
apply (*intro order.antisym inf-greatest*)
apply (*rule Inf-greatest, safe*)
apply (*case-tac n*)
apply *simp-all*
using *Inf-lower inf.coboundedI2* **apply** *force*
apply (*simp add: Inf-lower*)
by (*auto intro: Inf-mono*)

end

lemma *Sup-sup-closed*: *Sup-closed-set* $(X :: 'a :: complete-lattice \ set) \Longrightarrow \text{sup-closed-set } X$
by (*metis Sup-closed-set-def empty-subsetI insert-subsetI sup-Sup sup-closed-set-def*)

lemma *Inf-inf-closed*: *Inf-closed-set* $(X :: 'a :: complete-lattice \ set) \Longrightarrow \text{inf-closed-set } X$
by (*metis Inf-closed-set-def empty-subsetI inf-Inf inf-closed-set-def insert-subset*)

4.5 Sup- and Inf-Preservation

Next, important notation for morphism between posets and lattices is introduced: sup-preservation, inf-preservation and related properties.

abbreviation *Sup-pres* :: $('a :: Sup \Rightarrow 'b :: Sup) \Rightarrow \text{bool}$ **where**
Sup-pres $f \equiv f \circ Sup = Sup \circ (\cdot) f$

abbreviation $Inf-pres :: ('a::Inf \Rightarrow 'b::Inf) \Rightarrow bool$ **where**
 $Inf-pres\ f \equiv f \circ Inf = Inf \circ (\cdot)\ f$

abbreviation $sup-pres :: ('a::sup \Rightarrow 'b::sup) \Rightarrow bool$ **where**
 $sup-pres\ f \equiv (\forall x\ y. f\ (x \sqcup y) = f\ x \sqcup f\ y)$

abbreviation $inf-pres :: ('a::inf \Rightarrow 'b::inf) \Rightarrow bool$ **where**
 $inf-pres\ f \equiv (\forall x\ y. f\ (x \sqcap y) = f\ x \sqcap f\ y)$

abbreviation $bot-pres :: ('a::bot \Rightarrow 'b::bot) \Rightarrow bool$ **where**
 $bot-pres\ f \equiv f\ \perp = \perp$

abbreviation $top-pres :: ('a::top \Rightarrow 'b::top) \Rightarrow bool$ **where**
 $top-pres\ f \equiv f\ \top = \top$

abbreviation $Sup-dual :: ('a::Sup \Rightarrow 'b::Inf) \Rightarrow bool$ **where**
 $Sup-dual\ f \equiv f \circ Sup = Inf \circ (\cdot)\ f$

abbreviation $Inf-dual :: ('a::Inf \Rightarrow 'b::Sup) \Rightarrow bool$ **where**
 $Inf-dual\ f \equiv f \circ Inf = Sup \circ (\cdot)\ f$

abbreviation $sup-dual :: ('a::sup \Rightarrow 'b::inf) \Rightarrow bool$ **where**
 $sup-dual\ f \equiv (\forall x\ y. f\ (x \sqcup y) = f\ x \sqcap f\ y)$

abbreviation $inf-dual :: ('a::inf \Rightarrow 'b::sup) \Rightarrow bool$ **where**
 $inf-dual\ f \equiv (\forall x\ y. f\ (x \sqcap y) = f\ x \sqcup f\ y)$

abbreviation $bot-dual :: ('a::bot \Rightarrow 'b::top) \Rightarrow bool$ **where**
 $bot-dual\ f \equiv f\ \perp = \top$

abbreviation $top-dual :: ('a::top \Rightarrow 'b::bot) \Rightarrow bool$ **where**
 $top-dual\ f \equiv f\ \top = \perp$

Inf-preservation and sup-preservation relate with duality.

lemma $Inf-pres-map-dual-var$:

$Inf-pres\ f = Sup-pres\ (\partial_F\ f)$

for $f :: 'a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual$

proof –

{ **fix** $x :: 'a\ set$

assume $\partial\ (f\ (\sqcap\ (\partial\ 'x))) = (\sqcup_{y \in x}. \partial\ (f\ (\partial\ y)))$ **for** x

then have $\sqcap\ (f\ ' \partial\ 'A) = f\ (\partial\ (\sqcup\ A))$ **for** A

by ($metis\ (no-types)\ Sup-dual-def-var\ image-image\ invol-dual-var\ subset-dual$)

then have $\sqcap\ (f\ 'x) = f\ (\sqcap\ x)$

by ($metis\ Sup-dual-def-var\ subset-dual$) }

then show $?thesis$

by ($auto\ simp\ add:\ map-dual-def\ fun-eq-iff\ Inf-dual-var\ Sup-dual-def-var\ image-comp$)

qed

lemma *Inf-pres-map-dual*: $Inf\text{-pres} = Sup\text{-pres} \circ (\partial_F :: ('a :: complete\text{-lattice-with-dual} \Rightarrow 'b :: complete\text{-lattice-with-dual}) \Rightarrow 'a \Rightarrow 'b)$

proof –

```
{fix f :: 'a ⇒ 'b
  have Inf-pres f = (Sup-pres ∘ ∂F) f
    by (simp add: Inf-pres-map-dual-var)}
thus ?thesis
  by force
```

qed

lemma *Sup-pres-map-dual-var*:

```
fixes f :: 'a :: complete-lattice-with-dual ⇒ 'b :: complete-lattice-with-dual
shows Sup-pres f = Inf-pres (∂F f)
by (metis Inf-pres-map-dual-var fun-dual5 map-dual-def)
```

lemma *Sup-pres-map-dual*: $Sup\text{-pres} = Inf\text{-pres} \circ (\partial_F :: ('a :: complete\text{-lattice-with-dual} \Rightarrow 'b :: complete\text{-lattice-with-dual}) \Rightarrow 'a \Rightarrow 'b)$

by (*simp add: Inf-pres-map-dual comp-assoc map-dual-invol*)

The following lemmas relate isotonicity of functions between complete lattices with weak (left) preservation properties of sups and infs.

lemma *fun-isol*: $mono\ f \implies mono\ ((\circ)\ f)$

by (*simp add: le-fun-def mono-def*)

lemma *fun-isor*: $mono\ f \implies mono\ (\lambda x. x \circ f)$

by (*simp add: le-fun-def mono-def*)

lemma *Sup-sup-pres*:

```
fixes f :: 'a :: complete-lattice ⇒ 'b :: complete-lattice
shows Sup-pres f ⇒ sup-pres f
by (metis (no-types, opaque-lifting) Sup-empty Sup-insert comp-apply image-insert
sup-bot.right-neutral)
```

lemma *Inf-inf-pres*:

```
fixes f :: 'a :: complete-lattice ⇒ 'b :: complete-lattice
shows Inf-pres f ⇒ inf-pres f
by (smt (verit) INF-insert Inf-empty Inf-insert comp-eq-elim inf-top.right-neutral)
```

lemma *Sup-bot-pres*:

```
fixes f :: 'a :: complete-lattice ⇒ 'b :: complete-lattice
shows Sup-pres f ⇒ bot-pres f
by (metis SUP-empty Sup-empty comp-eq-elim)
```

lemma *Inf-top-pres*:

```
fixes f :: 'a :: complete-lattice ⇒ 'b :: complete-lattice
shows Inf-pres f ⇒ top-pres f
by (metis INF-empty Inf-empty comp-eq-elim)
```

lemma *Sup-sup-dual*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{Sup-dual } f \Longrightarrow \text{sup-dual } f$
by (*smt (verit) comp-eq-elim image-empty image-insert inf-Inf sup-Sup*)

lemma *Inf-inf-dual*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{Inf-dual } f \Longrightarrow \text{inf-dual } f$
by (*smt (verit) comp-eq-elim image-empty image-insert inf-Inf sup-Sup*)

lemma *Sup-bot-dual*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{Sup-dual } f \Longrightarrow \text{bot-dual } f$
by (*metis INF-empty Sup-empty comp-eq-elim*)

lemma *Inf-top-dual*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{Inf-dual } f \Longrightarrow \text{top-dual } f$
by (*metis Inf-empty SUP-empty comp-eq-elim*)

However, Inf-preservation does not imply top-preservation and Sup-preservation does not imply bottom-preservation.

lemma
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{Sup-pres } f \Longrightarrow \text{top-pres } f$
oops

lemma
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{Inf-pres } f \Longrightarrow \text{bot-pres } f$
oops

context *complete-lattice*
begin

lemma *iso-Inf-subdistl*:
fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$
shows $\text{mono } f \Longrightarrow f \circ \text{Inf} \leq \text{Inf} \circ (\cdot) f$
by (*simp add: complete-lattice-class.le-Inf-iff le-funI Inf-lower monoD*)

lemma *iso-Sup-supdistl*:
fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$
shows $\text{mono } f \Longrightarrow \text{Sup} \circ (\cdot) f \leq f \circ \text{Sup}$
by (*simp add: complete-lattice-class.Sup-le-iff le-funI Sup-upper monoD*)

lemma *Inf-subdistl-iso*:
fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$
shows $f \circ \text{Inf} \leq \text{Inf} \circ (\cdot) f \Longrightarrow \text{mono } f$
unfolding *mono-def le-fun-def comp-def* **by** (*metis complete-lattice-class.le-INF-iff*)

Inf-atLeast atLeast-iff)

lemma *Sup-supdistl-iso*:

fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$

shows $\text{Sup} \circ (\cdot) f \leq f \circ \text{Sup} \implies \text{mono } f$

unfolding *mono-def le-fun-def comp-def* **by** (*metis complete-lattice-class.SUP-le-iff Sup-atMost atMost-iff*)

lemma *supdistl-iso*:

fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$

shows $(\text{Sup} \circ (\cdot) f \leq f \circ \text{Sup}) = \text{mono } f$

using *Sup-supdistl-iso iso-Sup-supdistl* **by** *force*

lemma *subdistl-iso*:

fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$

shows $(f \circ \text{Inf} \leq \text{Inf} \circ (\cdot) f) = \text{mono } f$

using *Inf-subdistl-iso iso-Inf-subdistl* **by** *force*

end

lemma *ord-iso-Inf-pres*:

fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$

shows *ord-iso* $f \implies \text{Inf} \circ (\cdot) f = f \circ \text{Inf}$

proof –

let $?g = \text{the-inv } f$

assume $h: \text{ord-iso } f$

hence $a: \text{mono } ?g$

by (*simp add: ord-iso-the-inv*)

{fix $X :: 'a::\text{complete-lattice set}$

{fix $y :: 'b::\text{complete-lattice}$

have $(y \leq f (\bigsqcap X)) = (?g y \leq \bigsqcap X)$

by (*metis (mono-tags, lifting) UNIV-I f-the-inv-into-f h monoD ord-embed-alt ord-embed-inj ord-iso-alt*)

also have $\dots = (\forall x \in X. ?g y \leq x)$

by (*simp add: le-Inf-iff*)

also have $\dots = (\forall x \in X. y \leq f x)$

by (*metis (mono-tags, lifting) UNIV-I f-the-inv-into-f h monoD ord-embed-alt ord-embed-inj ord-iso-alt*)

also have $\dots = (y \leq \bigsqcap (f \text{ ` } X))$

by (*simp add: le-INF-iff*)

finally have $(y \leq f (\bigsqcap X)) = (y \leq \bigsqcap (f \text{ ` } X)).$

hence $f (\bigsqcap X) = \bigsqcap (f \text{ ` } X)$

by (*meson dual-order.antisym order-refl*)}

thus $?thesis$

unfolding *fun-eq-iff* **by** *simp*

qed

lemma *ord-iso-Sup-pres*:

fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$

shows $\text{ord-iso } f \implies \text{Sup} \circ (\cdot) f = f \circ \text{Sup}$
proof –
let $?g = \text{the-inv } f$
assume $h: \text{ord-iso } f$
hence $a: \text{mono } ?g$
by (*simp add: ord-iso-the-inv*)
{fix $X :: 'a::\text{complete-lattice set}$
{fix $y :: 'b::\text{complete-lattice}$
have $(f (\bigsqcup X) \leq y) = (\bigsqcup X \leq ?g y)$
by (*metis (mono-tags, lifting) UNIV-I f-the-inv-into-f h monoD ord-embed-alt ord-embed-inj ord-iso-alt*)
also have $\dots = (\forall x \in X. x \leq ?g y)$
by (*simp add: Sup-le-iff*)
also have $\dots = (\forall x \in X. f x \leq y)$
by (*metis (mono-tags, lifting) UNIV-I f-the-inv-into-f h monoD ord-embed-alt ord-embed-inj ord-iso-alt*)
also have $\dots = (\bigsqcup (f ' X) \leq y)$
by (*simp add: SUP-le-iff*)
finally have $(f (\bigsqcup X) \leq y) = (\bigsqcup (f ' X) \leq y).$
hence $f (\bigsqcup X) = \bigsqcup (f ' X)$
by (*meson dual-order.antisym order-refl*)
thus $?thesis$
unfolding *fun-eq-iff* **by** *simp*
qed

Right preservation of sups and infs is trivial.

lemma *fSup-distr: Sup-pres* $(\lambda x. x \circ f)$
unfolding *fun-eq-iff* **by** (*simp add: image-comp*)

lemma *fSup-distr-var: $\bigsqcup F \circ g = (\bigsqcup f \in F. f \circ g)$*
unfolding *fun-eq-iff* **by** (*simp add: image-comp*)

lemma *fInf-distr: Inf-pres* $(\lambda x. x \circ f)$
unfolding *fun-eq-iff comp-def*
by (*smt (verit) INF-apply Inf-fun-def Sup.SUP-cong*)

lemma *fInf-distr-var: $\prod F \circ g = (\prod f \in F. f \circ g)$*
unfolding *fun-eq-iff comp-def*
by (*smt (verit) INF-apply INF-cong INF-image Inf-apply image-comp image-def image-image*)

The next set of lemma revisits the preservation properties in the function space.

lemma *fSup-subdistl:*
assumes $\text{mono } (f::'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice})$
shows $\text{Sup} \circ (\cdot) ((\circ) f) \leq (\circ) f \circ \text{Sup}$
using *assms* **by** (*simp add: fun-isol supdistl-iso*)

lemma *fSup-subdistl-var:*

fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{mono } f \Longrightarrow (\bigsqcup g \in G. f \circ g) \leq f \circ \bigsqcup G$
by (*simp add: fun-isol mono-Sup*)

lemma *fInf-subdistl*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{mono } f \Longrightarrow (\circ) f \circ \text{Inf} \leq \text{Inf} \circ (\circ) ((\circ) f)$
by (*simp add: fun-isol subdistl-iso*)

lemma *fInf-subdistl-var*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{mono } f \Longrightarrow f \circ \prod G \leq (\prod g \in G. f \circ g)$
by (*simp add: fun-isol mono-Inf*)

lemma *fSup-distl*: $\text{Sup-pres } f \Longrightarrow \text{Sup-pres } ((\circ) f)$
unfolding *fun-eq-iff* **by** (*simp add: image-comp*)

lemma *fSup-distl-var*: $\text{Sup-pres } f \Longrightarrow f \circ \bigsqcup G = (\bigsqcup g \in G. f \circ g)$
unfolding *fun-eq-iff* **by** (*simp add: image-comp*)

lemma *fInf-distl*: $\text{Inf-pres } f \Longrightarrow \text{Inf-pres } ((\circ) f)$
unfolding *fun-eq-iff* **by** (*simp add: image-comp*)

lemma *fInf-distl-var*: $\text{Inf-pres } f \Longrightarrow f \circ \prod G = (\prod g \in G. f \circ g)$
unfolding *fun-eq-iff* **by** (*simp add: image-comp*)

Downsets preserve infs whereas upsets preserve sups.

lemma *Inf-pres-downset*: $\text{Inf-pres } (\downarrow :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a \text{ set})$
unfolding *downset-prop fun-eq-iff*
by (*safe, simp-all add: le-Inf-iff*)

lemma *Sup-dual-upset*: $\text{Sup-dual } (\uparrow :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a \text{ set})$
unfolding *upset-prop fun-eq-iff*
by (*safe, simp-all add: Sup-le-iff*)

Images of Sup-morphisms are closed under Sups and images of Inf-morphisms are closed under Infs.

lemma *Sup-pres-Sup-closed*: $\text{Sup-pres } f \Longrightarrow \text{Sup-closed-set } (\text{range } f)$
by (*metis (mono-tags, lifting) Sup-closed-set-def comp-eq-elim range-eqI subset-image-iff*)

lemma *Inf-pres-Inf-closed*: $\text{Inf-pres } f \Longrightarrow \text{Inf-closed-set } (\text{range } f)$
by (*metis (mono-tags, lifting) Inf-closed-set-def comp-eq-elim range-eqI subset-image-iff*)

It is well known that functions into complete lattices form complete lattices. Here, such results are shown for the subclasses of isotone functions, where additional closure conditions must be respected.

typedef (**overloaded**) $'a \text{ iso} = \{f :: 'a::\text{order} \Rightarrow 'a::\text{order}. \text{mono } f\}$

```

    by (metis Abs-ord-homset-cases ord-homset-def)

setup-lifting type-definition-iso

instantiation iso :: (complete-lattice) complete-lattice
begin

lift-definition Inf-iso :: 'a::complete-lattice iso set  $\Rightarrow$  'a iso is Sup
    by (metis (mono-tags, lifting) SUP-subset-mono Sup-apply mono-def subsetI)

lift-definition Sup-iso :: 'a::complete-lattice iso set  $\Rightarrow$  'a iso is Inf
    by (smt (verit) INF-lower2 Inf-apply le-INF-iff mono-def)

lift-definition bot-iso :: 'a::complete-lattice iso is  $\top$ 
    by (simp add: monoI)

lift-definition sup-iso :: 'a::complete-lattice iso  $\Rightarrow$  'a iso  $\Rightarrow$  'a iso is inf
    by (smt (verit) inf-apply inf-mono monoD monoI)

lift-definition top-iso :: 'a::complete-lattice iso is  $\perp$ 
    by (simp add: mono-def)

lift-definition inf-iso :: 'a::complete-lattice iso  $\Rightarrow$  'a iso  $\Rightarrow$  'a iso is sup
    by (smt (verit) mono-def sup.mono sup-apply)

lift-definition less-eq-iso :: 'a::complete-lattice iso  $\Rightarrow$  'a iso  $\Rightarrow$  bool is  $(\geq)$ .

lift-definition less-iso :: 'a::complete-lattice iso  $\Rightarrow$  'a iso  $\Rightarrow$  bool is  $(>)$ .

instance
    by (intro-classes; transfer, simp-all add: less-fun-def Sup-upper Sup-least Inf-lower
        Inf-greatest)

end

```

Duality has been baked into this result because of its relevance for predicate transformers. A proof where Sups are mapped to Sups and Infs to Infs is certainly possible, but two instantiation of the same type and the same classes are unfortunately impossible. Interpretations could be used instead. A corresponding result for Inf-preserving functions and Sup-lattices, is proved in components on transformers, as more advanced properties about Inf-preserving functions are needed.

4.6 Alternative Definitions for Complete Boolean Algebras

The current definitions of complete boolean algebras deviates from that in most textbooks in that a distributive law with infinite sups and infinite infs is used. There are interesting applications, for instance in topology, where

weaker laws are needed — for instance for frames and locales.

```
class complete-heyting-algebra = complete-lattice +
  assumes ch-dist:  $x \sqcap \bigsqcup Y = (\bigsqcup y \in Y. x \sqcap y)$ 
```

Complete Heyting algebras are also known as frames or locales (they differ with respect to their morphisms).

```
class complete-co-heyting-algebra = complete-lattice +
  assumes co-ch-dist:  $x \sqcup \bigsqcap Y = (\bigsqcap y \in Y. x \sqcup y)$ 
```

```
class complete-boolean-algebra-alt = complete-lattice + boolean-algebra
```

```
instance set :: (type) complete-boolean-algebra-alt..
```

```
context complete-boolean-algebra-alt
begin
```

```
subclass complete-heyting-algebra
```

```
proof
```

```
  fix x Y
```

```
  {fix t
```

```
    have ( $x \sqcap \bigsqcup Y \leq t$ ) = ( $\bigsqcup Y \leq -x \sqcup t$ )
```

```
      by (simp add: inf commute shunt1[symmetric])
```

```
    also have ... = ( $\forall y \in Y. y \leq -x \sqcup t$ )
```

```
      using Sup-le-iff by blast
```

```
    also have ... = ( $\forall y \in Y. x \sqcap y \leq t$ )
```

```
      by (simp add: inf commute shunt1)
```

```
    finally have ( $x \sqcap \bigsqcup Y \leq t$ ) = ( $(\bigsqcup y \in Y. x \sqcap y) \leq t$ )
```

```
      by (simp add: local.SUP-le-iff)}
```

```
  thus  $x \sqcap \bigsqcup Y = (\bigsqcup y \in Y. x \sqcap y)$ 
```

```
    using order.eq-iff by blast
```

```
qed
```

```
subclass complete-co-heyting-algebra
```

```
  apply unfold-locales
```

```
  apply (rule order.antisym)
```

```
  apply (simp add: INF-greatest Inf-lower2)
```

```
  by (meson eq-refl le-INF-iff le-Inf-iff shunt2)
```

```
lemma de-morgan1:  $-(\bigsqcup X) = (\bigsqcap x \in X. -x)$ 
```

```
proof -
```

```
  {fix y
```

```
    have ( $y \leq -(\bigsqcup X)$ ) = ( $\bigsqcup X \leq -y$ )
```

```
      using compl-le-swap1 by blast
```

```
    also have ... = ( $\forall x \in X. x \leq -y$ )
```

```
      by (simp add: Sup-le-iff)
```

```
    also have ... = ( $\forall x \in X. y \leq -x$ )
```

```
      using compl-le-swap1 by blast
```

```
    also have ... = ( $y \leq (\bigsqcap x \in X. -x)$ )
```

```
      using le-INF-iff by force
```

```

finally have  $(y \leq -(\bigsqcup X)) = (y \leq (\bigsqcap x \in X. -x)).\}$ 
thus ?thesis
  using order.antisym by blast
qed

lemma de-morgan2:  $-(\bigsqcap X) = (\bigsqcup x \in X. -x)$ 
  by (metis de-morgan1 ba-dual.dual-iff ba-dual.image-dual pointfree-idE)

end

class complete-boolean-algebra-alt-with-dual = complete-lattice-with-dual + complete-boolean-algebra-alt

instantiation set :: (type) complete-boolean-algebra-alt-with-dual
begin

definition dual-set :: 'a set  $\Rightarrow$  'a set where
  dual-set = uminus

instance
  by intro-classes (simp-all add: ba-dual.inj-dual dual-set-def comp-def uminus-Sup id-def)

end

context complete-boolean-algebra-alt
begin

sublocale cba-dual: complete-boolean-algebra-alt-with-dual - - - - - uminus - -
  by unfold-locales (auto simp: de-morgan2 de-morgan1)

end

```

4.7 Atomic Boolean Algebras

Next, atomic boolean algebras are defined.

```
context bounded-lattice
```

```
begin
```

Atoms are covers of bottom.

```
definition atom  $x = (x \neq \perp \wedge \neg(\exists y. \perp < y \wedge y < x))$ 
```

```
definition atom-map  $x = \{y. \text{atom } y \wedge y \leq x\}$ 
```

```
lemma atom-map-def-var: atom-map  $x = \downarrow x \cap \text{Collect } \text{atom}$ 
```

```
  unfolding atom-map-def downset-def downset-set-def comp-def atom-def by fast-force
```

```
lemma atom-map-atoms:  $\bigcup (\text{range } \text{atom-map}) = \text{Collect } \text{atom}$ 
```

```

unfolding atom-map-def atom-def by auto

end

typedef (overloaded) 'a atoms = range (atom-map::'a::bounded-lattice ⇒ 'a set)
by blast

setup-lifting type-definition-atoms

definition at-map :: 'a::bounded-lattice ⇒ 'a atoms where
  at-map = Abs-atoms ◦ atom-map

class atomic-boolean-algebra = boolean-algebra +
  assumes atomicity: x ≠ ⊥ ⇒ (∃ y. atom y ∧ y ≤ x)

class complete-atomic-boolean-algebra = complete-lattice + atomic-boolean-algebra

begin

subclass complete-boolean-algebra-alt..

end

```

Here are two equivalent definitions for atoms; first in boolean algebras, and then in complete boolean algebras.

```

context boolean-algebra
begin

```

The following two conditions are taken from Koppelberg's book [6].

```

lemma atom-neg: atom x ⇒ x ≠ ⊥ ∧ (∀ y z. x ≤ y ∨ x ≤ -y)
by (auto simp add: atom-def) (metis local.dual-order.not-eq-order-implies-strict
local.inf.cobounded1 local.inf.cobounded2 local.inf-shunt)

```

```

lemma atom-sup: (∀ y. x ≤ y ∨ x ≤ -y) ⇒ (∀ y z. (x ≤ y ∨ x ≤ z) = (x ≤ y ⊔ z))
by (metis inf.orderE le-supI1 shunt2)

```

```

lemma sup-atom: x ≠ ⊥ ⇒ (∀ y z. (x ≤ y ∨ x ≤ z) = (x ≤ y ⊔ z)) ⇒ atom x
by (auto simp add: atom-def) (metis (full-types) local.inf.boundedI local.inf.cobounded2
local.inf-shunt local.inf-sup-ord(4) local.le-iff-sup local.shunt1 local.sup.absorb1 local.sup.strict-order-iff)

```

```

lemma atom-sup-iff: atom x = (x ≠ ⊥ ∧ (∀ y z. (x ≤ y ∨ x ≤ z) = (x ≤ y ⊔ z)))
by rule (auto simp add: atom-neg atom-sup sup-atom)

```

```

lemma atom-neg-iff: atom x = (x ≠ ⊥ ∧ (∀ y z. x ≤ y ∨ x ≤ -y))
by rule (auto simp add: atom-neg atom-sup sup-atom)

```

```

lemma atom-map-bot-pres: atom-map ⊥ = {}

```

```

using atom-def atom-map-def le-bot by auto

lemma atom-map-top-pres: atom-map  $\top = \text{Collect atom}$ 
  using atom-map-def by auto

end

context complete-boolean-algebra-alt
begin

lemma atom-Sup:  $\bigwedge Y. x \neq \perp \implies (\forall y. x \leq y \vee x \leq -y) \implies ((\exists y \in Y. x \leq y) = (x \leq \bigsqcup Y))$ 
  by (metis Sup-least Sup-upper2 compl-le-swap1 le-iff-inf inf-shunt)

lemma Sup-atom:  $x \neq \perp \implies (\forall Y. (\exists y \in Y. x \leq y) = (x \leq \bigsqcup Y)) \implies \text{atom } x$ 
proof -
  assume h1:  $x \neq \perp$ 
  and h2:  $\forall Y. (\exists y \in Y. x \leq y) = (x \leq \bigsqcup Y)$ 
  hence  $\forall y z. (x \leq y \vee x \leq z) = (x \leq y \sqcup z)$ 
  by (smt (verit) insert-iff sup-Sup sup-bot.right-neutral)
  thus atom x
  by (simp add: h1 sup-atom)
qed

lemma atom-Sup-iff:  $\text{atom } x = (x \neq \perp \wedge (\forall Y. (\exists y \in Y. x \leq y) = (x \leq \bigsqcup Y)))$ 
  by standard (auto simp: atom-neg atom-Sup Sup-atom)

end

end

```

5 Representation Theorems for Orderings and Lattices

```

theory Representations
  imports Order-Lattice-Props

```

```
begin
```

5.1 Representation of Posets

The isomorphism between partial orders and downsets with set inclusion is well known. It forms the basis of Priestley and Stone duality. I show it not only for objects, but also order morphisms, hence establish equivalences and isomorphisms between categories.

```

typedef (overloaded) 'a downset = range ( $\downarrow :: 'a :: \text{ord} \Rightarrow 'a \text{ set}$ )
  by fastforce

```

setup-lifting *type-definition-downset*

The map ds yields the isomorphism between the set and the powerset level if its range is restricted to downsets.

definition $ds :: 'a::ord \Rightarrow 'a \text{ \textit{downset}}$ **where**
 $ds = Abs\text{-downset} \circ \downarrow$

In a complete lattice, its inverse is Sup .

definition $SSup :: 'a::complete\text{-lattice} \text{ \textit{downset}} \Rightarrow 'a$ **where**
 $SSup = Sup \circ Rep\text{-downset}$

lemma $ds\text{-}SSup\text{-inv}$: $ds \circ SSup = (id :: 'a::complete\text{-lattice} \text{ \textit{downset}} \Rightarrow 'a \text{ \textit{downset}})$
unfolding $ds\text{-def}$ $SSup\text{-def}$
by (*smt* (*verit*) *Rep-downset* *Rep-downset-inverse* *cSup-atMost* *eq-id-iff* *imageE* *o-def* *ord-class.atMost-def* *ord-class.downset-prop*)

lemma $SSup\text{-}ds\text{-inv}$: $SSup \circ ds = (id :: 'a::complete\text{-lattice} \Rightarrow 'a)$
unfolding $ds\text{-def}$ $SSup\text{-def}$ *fun-eq-iff* *id-def* *comp-def* **by** (*simp* *add*: *Abs-downset-inverse* *pointfree-idE*)

instantiation $downset :: (ord) \text{ \textit{order}}$
begin

lift-definition $less\text{-eq}\text{-downset} :: 'a \text{ \textit{downset}} \Rightarrow 'a \text{ \textit{downset}} \Rightarrow bool$ **is** $(\lambda X Y. Rep\text{-downset } X \subseteq Rep\text{-downset } Y)$.

lift-definition $less\text{-downset} :: 'a \text{ \textit{downset}} \Rightarrow 'a \text{ \textit{downset}} \Rightarrow bool$ **is** $(\lambda X Y. Rep\text{-downset } X \subset Rep\text{-downset } Y)$.

instance
by (*intro-classes*, *transfer*, *auto* *simp*: *Rep-downset-inject* *less-eq-downset-def*)

end

lemma $ds\text{-iso}$: *mono* ds
unfolding *mono-def* $ds\text{-def}$ *fun-eq-iff* *comp-def*
by (*metis* *Abs-downset-inverse* *downset-iso-iff* *less-eq-downset.rep-eq* *rangeI*)

lemma $ds\text{-faithful}$: $ds \ x \leq ds \ y \Longrightarrow x \leq (y :: 'a::order)$
by (*simp* *add*: *Abs-downset-inverse* *downset-faithful* $ds\text{-def}$ *less-eq-downset.rep-eq*)

lemma $ds\text{-inj}$: *inj* $(ds :: 'a::order \Rightarrow 'a \text{ \textit{downset}})$
by (*simp* *add*: $ds\text{-faithful}$ *dual-order.antisym* *injI*)

lemma $ds\text{-surj}$: *surj* ds
by (*metis* (*no-types*, *opaque-lifting*) *Rep-downset* *Rep-downset-inverse* $ds\text{-def}$ *image-iff* *o-apply* *surj-def*)

lemma *ds-bij*: *bij* (*ds*::*'a*::*order* \Rightarrow *'a* *downset*)
by (*simp add*: *bijI ds-inj ds-surj*)

lemma *ds-ord-iso*: *ord-iso ds*
unfolding *ord-iso-def comp-def inf-bool-def* **by** (*smt (verit) UNIV-I ds-bij ds-faithful ds-inj ds-iso ds-surj f-the-inv-into-f infI mono-def*)

The morphisms between orderings and downsets are isotone functions. One can define functors mapping back and forth between these.

definition *map-ds* :: (*'a*::*complete-lattice* \Rightarrow *'b*::*complete-lattice*) \Rightarrow (*'a* *downset* \Rightarrow *'b* *downset*) **where**
map-ds f = *ds* \circ *f* \circ *SSup*

This definition is actually contrived. We have shown that a function *f* between posets *P* and *Q* is isotone if and only if the inverse image of *f* maps downclosed sets in *Q* to downclosed sets in *P*. There is the following duality: *ds* is a natural transformation between the identity functor and the preimage functor as a contravariant functor from *P* to *Q*. Hence orderings with isotone maps and downsets with downset-preserving maps are dual, which is a first step towards Stone duality. I don't see a way of proving this with Isabelle, as the types of the preimage of *f* are the wrong way and I don't see how I could capture opposition with what I have.

lemma *map-ds-prop*:
fixes *f* :: *'a*::*complete-lattice* \Rightarrow *'b*::*complete-lattice*
shows *map-ds f* \circ *ds* = *ds* \circ *f*
unfolding *map-ds-def* **by** (*simp add*: *SSup-ds-inv comp-assoc*)

lemma *map-ds-prop2*:
fixes *f* :: *'a*::*complete-lattice* \Rightarrow *'b*::*complete-lattice*
shows *map-ds f* \circ *ds* = *ds* \circ *id f*
unfolding *map-ds-def* **by** (*simp add*: *SSup-ds-inv comp-assoc*)

This is part of showing that *map-ds* is naturally isomorphic to the identity functor, *ds* being the natural isomorphism.

definition *map-SSup* :: (*'a* *downset* \Rightarrow *'b* *downset*) \Rightarrow (*'a*::*complete-lattice* \Rightarrow *'b*::*complete-lattice*) **where**
map-SSup F = *SSup* \circ *F* \circ *ds*

lemma *map-ds-iso-pres*:
fixes *f* :: *'a*::*complete-lattice* \Rightarrow *'b*::*complete-lattice*
shows *mono f* \implies *mono (map-ds f)*
unfolding *fun-eq-iff mono-def map-ds-def ds-def SSup-def comp-def*
by (*metis Abs-downset-inverse Sup-subset-mono downset-iso-iff less-eq-downset.rep-eq rangeI*)

lemma *map-SSup-iso-pres*:
fixes *F* :: *'a*::*complete-lattice* *downset* \Rightarrow *'b*::*complete-lattice* *downset*

shows $\text{mono } F \implies \text{mono } (\text{map-SSup } F)$
unfolding $\text{fun-eq-iff mono-def map-SSup-def ds-def SSup-def comp-def}$
by ($\text{metis Abs-downset-inverse Sup-subset-mono downset-iso-iff less-eq-downset.rep-eq rangeI}$)

lemma map-SSup-prop :

fixes $F :: 'a::\text{complete-lattice downset} \Rightarrow 'b::\text{complete-lattice downset}$
shows $ds \circ \text{map-SSup } F = F \circ ds$
unfolding map-SSup-def **by** ($\text{metis ds-SSup-inv fun.map-id0 id-def rewriteL-comp-comp}$)

lemma map-SSup-prop2 :

fixes $F :: 'a::\text{complete-lattice downset} \Rightarrow 'b::\text{complete-lattice downset}$
shows $ds \circ \text{map-SSup } F = \text{id } F \circ ds$
by ($\text{simp add: map-SSup-prop}$)

lemma map-ds-func1 : $\text{map-ds id} = (\text{id}::'a::\text{complete-lattice downset} \Rightarrow 'a \text{ downset})$

by ($\text{simp add: ds-SSup-inv map-ds-def}$)

lemma map-ds-func2 :

fixes $g :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{map-ds } (f \circ g) = \text{map-ds } f \circ \text{map-ds } g$
unfolding $\text{map-ds-def fun-eq-iff comp-def ds-def SSup-def}$
by ($\text{metis Abs-downset-inverse Sup-atMost atMost-def downset-prop rangeI}$)

lemma map-SSup-func1 : $\text{map-SSup } (\text{id}::'a::\text{complete-lattice downset} \Rightarrow 'a \text{ downset}) = \text{id}$

by ($\text{simp add: SSup-ds-inv map-SSup-def}$)

lemma map-SSup-func2 :

fixes $F :: 'c::\text{complete-lattice downset} \Rightarrow 'b::\text{complete-lattice downset}$
and $G :: 'a::\text{complete-lattice downset} \Rightarrow 'c \text{ downset}$
shows $\text{map-SSup } (F \circ G) = \text{map-SSup } F \circ \text{map-SSup } G$
unfolding $\text{map-SSup-def fun-eq-iff comp-def id-def ds-def}$
by ($\text{metis comp-apply ds-SSup-inv ds-def id-apply}$)

lemma $\text{map-SSup-map-ds-inv}$: $\text{map-SSup} \circ \text{map-ds} = (\text{id}::('a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}) \Rightarrow ('a \Rightarrow 'b))$

by ($\text{metis (no-types, opaque-lifting) SSup-ds-inv comp-def eq-id-iff fun.map-comp fun.map-id0 id-apply map-SSup-prop map-ds-prop}$)

lemma $\text{map-ds-map-SSup-inv}$: $\text{map-ds} \circ \text{map-SSup} = (\text{id}::('a::\text{complete-lattice downset} \Rightarrow 'b::\text{complete-lattice downset}) \Rightarrow ('a \text{ downset} \Rightarrow 'b \text{ downset}))$

unfolding $\text{map-SSup-def map-ds-def SSup-def ds-def id-def comp-def fun-eq-iff}$
by ($\text{metis (no-types, lifting) Rep-downset Rep-downset-inverse Sup-downset-id image-iff pointfree-idE}$)

lemma inj-map-ds : $\text{inj } (\text{map-ds}::('a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}) \Rightarrow ('a \text{ downset} \Rightarrow 'b \text{ downset}))$

by ($\text{metis (no-types, lifting) SSup-ds-inv fun.map-id0 id-comp inj-def map-ds-prop}$)

rewriteR-comp-comp2)

lemma *inj-map-SSup*: *inj (map-SSup::('a::complete-lattice downset \Rightarrow 'b::complete-lattice downset) \Rightarrow ('a \Rightarrow 'b))*
by (*metis inj-on-id inj-on-imageI2 map-ds-map-SSup-inv*)

lemma *map-ds-map-SSup-iff*:
fixes *g :: 'a::complete-lattice \Rightarrow 'b::complete-lattice*
shows (*f = map-ds g*) = (*map-SSup f = g*)
by (*metis inj-eq inj-map-ds map-ds-map-SSup-inv pointfree-idE*)

This gives an isomorphism between categories.

lemma *surj-map-ds*: *surj (map-ds::('a::complete-lattice \Rightarrow 'b::complete-lattice) \Rightarrow ('a downset \Rightarrow 'b downset))*
by (*simp add: map-ds-map-SSup-iff surj-def*)

lemma *surj-map-SSup*: *surj (map-SSup::('a::complete-lattice-with-dual downset \Rightarrow 'b::complete-lattice-with-dual downset) \Rightarrow ('a \Rightarrow 'b))*
by (*metis map-ds-map-SSup-iff surjI*)

There is of course a dual result for upsets with the reverse inclusion ordering. Once again, it seems impossible to capture the "real" duality that uses the inverse image functor.

typedef (**overloaded**) *'a upset* = *range (\uparrow ::'a::ord \Rightarrow 'a set)*
by *fastforce*

setup-lifting *type-definition-upset*

definition *us :: 'a::ord \Rightarrow 'a upset* **where**
us = Abs-upset \circ \uparrow

definition *IInf :: 'a::complete-lattice upset \Rightarrow 'a* **where**
IInf = Inf \circ Rep-upset

lemma *us-ds*: *us = Abs-upset \circ (\wedge) ∂ \circ Rep-downset \circ ds \circ (∂ ::'a::ord-with-dual \Rightarrow 'a)*

unfolding *us-def ds-def fun-eq-iff comp-def* **by** (*simp add: Abs-downset-inverse upset-to-downset2*)

lemma *IInf-SSup*: *IInf = ∂ \circ SSup \circ Abs-downset \circ (\wedge) (∂ ::'a::complete-lattice-with-dual \Rightarrow 'a) \circ Rep-upset*

unfolding *IInf-def SSup-def fun-eq-iff comp-def*
by (*metis (no-types, opaque-lifting) Abs-downset-inverse Rep-upset Sup-dual-def-var image-iff rangeI subset-dual upset-to-downset3*)

lemma *us-IInf-inv*: *us \circ IInf = (id::'a::complete-lattice-with-dual upset \Rightarrow 'a upset)*

unfolding *us-def IInf-def fun-eq-iff id-def comp-def*

by (*metis (no-types, lifting) Inf-upset-id Rep-upset Rep-upset-inverse f-the-inv-into-f pointfree-idE upset-inj*)

lemma *IInf-us-inv*: $IInf \circ us = (id :: 'a :: complete-lattice-with-dual \Rightarrow 'a)$

unfolding *us-def IInf-def fun-eq-iff id-def comp-def*

by (*metis Abs-upset-inverse Sup-Inf-var Sup-atLeastAtMost Sup-dual-upset-var order-refl rangeI*)

instantiation *upset* :: (*ord*) *order*

begin

lift-definition *less-eq-upset* :: '*a* *upset* \Rightarrow '*a* *upset* \Rightarrow *bool* **is** ($\lambda X Y. Rep-upset X \supseteq Rep-upset Y$).

lift-definition *less-upset* :: '*a* *upset* \Rightarrow '*a* *upset* \Rightarrow *bool* **is** ($\lambda X Y. Rep-upset X \supset Rep-upset Y$).

instance

by (*intro-classes, transfer, simp-all add: less-le-not-le less-eq-upset.rep-eq Rep-upset-inject*)

end

lemma *us-iso*: $x \leq y \Longrightarrow us\ x \leq us\ y$ (*us :: 'a :: order-with-dual*)

by (*simp add: Abs-upset-inverse less-eq-upset.rep-eq upset-anti-iff us-def*)

lemma *us-faithful*: $us\ x \leq us\ y \Longrightarrow x \leq y$ (*us :: 'a :: order-with-dual*)

by (*simp add: Abs-upset-inverse upset-faithful us-def less-eq-upset.rep-eq*)

lemma *us-inj*: *inj* (*us :: 'a :: order-with-dual* \Rightarrow '*a* *upset*)

unfolding *inj-def* **by** (*simp add: us-faithful dual-order.antisym*)

lemma *us-surj*: *surj* (*us :: 'a :: order-with-dual* \Rightarrow '*a* *upset*)

unfolding *surj-def* **by** (*metis Rep-upset Rep-upset-inverse comp-def image-iff us-def*)

lemma *us-bij*: *bij* (*us :: 'a :: order-with-dual* \Rightarrow '*a* *upset*)

by (*simp add: bij-def us-surj us-inj*)

lemma *us-ord-iso*: *ord-iso* (*us :: 'a :: order-with-dual* \Rightarrow '*a* *upset*)

unfolding *ord-iso-def*

by (*simp, metis (no-types, lifting) UNIV-I f-the-inv-into-f monoI us-iso us-bij us-faithful us-inj us-surj*)

definition *map-us* :: ('*a* :: *complete-lattice* \Rightarrow '*b* :: *complete-lattice*) \Rightarrow ('*a* *upset* \Rightarrow '*b* *upset*) **where**

$map-us\ f = us \circ f \circ IInf$

lemma *map-us-prop*: $map-us\ f \circ (us :: 'a :: complete-lattice-with-dual \Rightarrow 'a\ upset) = us \circ id\ f$

unfolding *map-us-def* **by** (*simp add: Inf-us-inv comp-assoc*)

definition *map-Inf* :: ('a upset \Rightarrow 'b upset) \Rightarrow ('a::complete-lattice \Rightarrow 'b::complete-lattice)

where

map-Inf *F* = *Inf* \circ *F* \circ *us*

lemma *map-Inf-prop*: (*us*::'a::complete-lattice-with-dual \Rightarrow 'a upset) \circ *map-Inf*
F = *id* *F* \circ *us*

proof –

have *us* \circ *map-Inf* *F* = (*us* \circ *Inf*) \circ *F* \circ *us*

by (*simp add: fun.map-comp map-Inf-def*)

thus *?thesis*

by (*simp add: us-Inf-inv*)

qed

lemma *map-us-func1*: *map-us id* = (*id*::'a::complete-lattice-with-dual upset \Rightarrow 'a upset)

unfolding *map-us-def fun-eq-iff comp-def us-def id-def Inf-def*

by (*metis (no-types, lifting) Inf-upset-id Rep-upset Rep-upset-inverse image-iff pointfree-idE*)

lemma *map-us-func2*:

fixes *f* :: 'c::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual

and *g* :: 'a::complete-lattice-with-dual \Rightarrow 'c

shows *map-us* (*f* \circ *g*) = *map-us* *f* \circ *map-us* *g*

unfolding *map-us-def fun-eq-iff comp-def us-def Inf-def*

by (*metis Abs-upset-inverse Sup-Inf-var Sup-atLeastAtMost Sup-dual-upset-var order-refl rangeI*)

lemma *map-Inf-func1*: *map-Inf id* = (*id*::'a::complete-lattice-with-dual \Rightarrow 'a)

unfolding *map-Inf-def fun-eq-iff comp-def id-def us-def Inf-def* **by** (*simp add: Abs-upset-inverse pointfree-idE*)

lemma *map-Inf-func2*:

fixes *F* :: 'c::complete-lattice-with-dual upset \Rightarrow 'b::complete-lattice-with-dual upset

and *G* :: 'a::complete-lattice-with-dual upset \Rightarrow 'c upset

shows *map-Inf* (*F* \circ *G*) = *map-Inf* *F* \circ *map-Inf* *G*

unfolding *map-Inf-def fun-eq-iff comp-def id-def us-def*

by (*metis comp-apply id-apply us-Inf-inv us-def*)

lemma *map-Inf-map-us-inv*: *map-Inf* \circ *map-us* = (*id*::('a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual) \Rightarrow ('a \Rightarrow 'b))

unfolding *map-Inf-def map-us-def Inf-def us-def id-def comp-def fun-eq-iff* **by**
(*simp add: Abs-upset-inverse pointfree-idE*)

lemma *map-us-map-Inf-inv*: *map-us* \circ *map-Inf* = (*id*::('a::complete-lattice-with-dual upset \Rightarrow 'b::complete-lattice-with-dual upset) \Rightarrow ('a upset \Rightarrow 'b upset))

unfolding *map-Inf-def map-us-def Inf-def us-def id-def comp-def fun-eq-iff*

by (*metis* (*no-types*, *lifting*) *Inf-upset-id* *Rep-upset* *Rep-upset-inverse* *image-iff* *pointfree-idE*)

lemma *inj-map-us*: *inj* (*map-us*::('a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual) \Rightarrow ('a *upset* \Rightarrow 'b *upset*))

unfolding *map-us-def* *us-def* *IInf-def* *inj-def* *comp-def* *fun-eq-iff*

by (*metis* (*no-types*, *opaque-lifting*) *Abs-upset-inverse* *Inf-upset-id* *pointfree-idE* *rangeI*)

lemma *inj-map-IInf*: *inj* (*map-IInf*::('a::complete-lattice-with-dual *upset* \Rightarrow 'b::complete-lattice-with-dual *upset*) \Rightarrow ('a \Rightarrow 'b))

unfolding *map-IInf-def* *fun-eq-iff* *inj-def* *comp-def* *IInf-def* *us-def*

by (*metis* (*no-types*, *opaque-lifting*) *Inf-upset-id* *Rep-upset* *Rep-upset-inverse* *image-iff* *pointfree-idE*)

lemma *map-us-map-IInf-iff*:

fixes *g* :: 'a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual

shows (*f* = *map-us* *g*) = (*map-IInf* *f* = *g*)

by (*metis* *inj-eq* *inj-map-us* *map-us-map-IInf-inv* *pointfree-idE*)

lemma *map-us-mono-pres*:

fixes *f* :: 'a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual

shows *mono* *f* \Longrightarrow *mono* (*map-us* *f*)

unfolding *mono-def* *map-us-def* *comp-def* *us-def* *IInf-def*

by (*metis* *Abs-upset-inverse* *Inf-superset-mono* *less-eq-upset.rep-eq* *rangeI* *upset-anti-iff*)

lemma *map-IInf-mono-pres*:

fixes *F* :: 'a::complete-lattice-with-dual *upset* \Rightarrow 'b::complete-lattice-with-dual *upset*

shows *mono* *F* \Longrightarrow *mono* (*map-IInf* *F*)

unfolding *mono-def* *map-IInf-def* *comp-def* *us-def* *IInf-def*

oops

lemma *surj-map-us*: *surj* (*map-us*::('a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual) \Rightarrow ('a *upset* \Rightarrow 'b *upset*))

by (*simp* *add*: *map-us-map-IInf-iff* *surj-def*)

lemma *surj-map-IInf*: *surj* (*map-IInf*::('a::complete-lattice-with-dual *upset* \Rightarrow 'b::complete-lattice-with-dual *upset*) \Rightarrow ('a \Rightarrow 'b))

by (*metis* *map-us-map-IInf-iff* *surjI*)

Hence we have again an isomorphism — or rather equivalence — between categories. Here, however, duality is not consistently picked up.

5.2 Stone's Theorem in the Presence of Atoms

Atom-map is a boolean algebra morphism.

context *boolean-algebra*

begin

lemma *atom-map-compl-pres*: $\text{atom-map } (-x) = \text{Collect atom} - \text{atom-map } x$

proof –

{**fix** y

have $(y \in \text{atom-map } (-x)) = (\text{atom } y \wedge y \leq -x)$

by (*simp add: atom-map-def*)

also have $\dots = (\text{atom } y \wedge \neg(y \leq x))$

by (*metis atom-sup-iff inf.orderE inf-shunt sup-compl-top top.ordering-top-axioms ordering-top.extremum*)

also have $\dots = (y \in \text{Collect atom} - \text{atom-map } x)$

using *atom-map-def* **by** *auto*

finally have $(y \in \text{atom-map } (-x)) = (y \in \text{Collect atom} - \text{atom-map } x).$

thus *?thesis*

by *blast*

qed

lemma *atom-map-sup-pres*: $\text{atom-map } (x \sqcup y) = \text{atom-map } x \cup \text{atom-map } y$

proof –

{**fix** z

have $(z \in \text{atom-map } (x \sqcup y)) = (\text{atom } z \wedge z \leq x \sqcup y)$

by (*simp add: atom-map-def*)

also have $\dots = (\text{atom } z \wedge (z \leq x \vee z \leq y))$

using *atom-sup-iff* **by** *auto*

also have $\dots = (z \in \text{atom-map } x \cup \text{atom-map } y)$

using *atom-map-def* **by** *auto*

finally have $(z \in \text{atom-map } (x \sqcup y)) = (z \in \text{atom-map } x \cup \text{atom-map } y)$

by *blast*

thus *?thesis*

by *blast*

qed

lemma *atom-map-inf-pres*: $\text{atom-map } (x \sqcap y) = \text{atom-map } x \cap \text{atom-map } y$

by (*smt (verit) Diff-Un atom-map-compl-pres atom-map-sup-pres compl-inf double-compl*)

lemma *atom-map-minus-pres*: $\text{atom-map } (x - y) = \text{atom-map } x - \text{atom-map } y$

using *atom-map-compl-pres atom-map-def diff-eq* **by** *auto*

end

The homomorphic images of boolean algebras under *atom-map* are boolean algebras — in fact powerset boolean algebras.

instantiation *atoms* :: (*boolean-algebra*) *boolean-algebra*

begin

lift-definition *minus-atoms* :: '*a atoms* \Rightarrow '*a atoms* \Rightarrow '*a atoms* **is** $\lambda x y. \text{Abs-atoms } (\text{Rep-atoms } x - \text{Rep-atoms } y).$

lift-definition *uminus-atoms* :: 'a atoms \Rightarrow 'a atoms **is** $\lambda x.$ Abs-atoms (Collect atom - Rep-atoms x).

lift-definition *bot-atoms* :: 'a atoms **is** Abs-atoms {}.

lift-definition *sup-atoms* :: 'a atoms \Rightarrow 'a atoms \Rightarrow 'a atoms **is** $\lambda x y.$ Abs-atoms (Rep-atoms x \cup Rep-atoms y).

lift-definition *top-atoms* :: 'a atoms **is** Abs-atoms (Collect atom).

lift-definition *inf-atoms* :: 'a atoms \Rightarrow 'a atoms \Rightarrow 'a atoms **is** $\lambda x y.$ Abs-atoms (Rep-atoms x \cap Rep-atoms y).

lift-definition *less-eq-atoms* :: 'a atoms \Rightarrow 'a atoms \Rightarrow bool **is** ($\lambda x y.$ Rep-atoms x \subseteq Rep-atoms y).

lift-definition *less-atoms* :: 'a atoms \Rightarrow 'a atoms \Rightarrow bool **is** ($\lambda x y.$ Rep-atoms x \subset Rep-atoms y).

instance

apply *intro-classes*

apply (transfer, simp add: less-le-not-le)

apply (transfer, simp)

apply (transfer, blast)

apply (simp add: Rep-atoms-inject less-eq-atoms.abs-eq)

apply (transfer, smt (verit) Abs-atoms-inverse Rep-atoms atom-map-inf-pres image-iff inf-le1 rangeI)

apply (transfer, smt (verit) Abs-atoms-inverse Rep-atoms atom-map-inf-pres image-iff inf-le2 rangeI)

apply (transfer, smt (verit) Abs-atoms-inverse Rep-atoms atom-map-inf-pres image-iff le-iff-sup rangeI sup-inf-distrib1)

apply (transfer, smt (verit) Abs-atoms-inverse Rep-atoms atom-map-sup-pres image-iff image-iff inf.orderE inf-sup-aci(6) le-iff-sup order-refl rangeI rangeI)

apply (transfer, smt (verit) Abs-atoms-inverse Rep-atoms atom-map-sup-pres image-iff inf-sup-aci(6) le-iff-sup rangeI sup.left-commute sup.right-idem)

apply (transfer, subst Abs-atoms-inverse, metis (no-types, lifting) Rep-atoms atom-map-sup-pres image-iff rangeI, simp)

apply transfer **using** Abs-atoms-inverse atom-map-bot-pres **apply** blast

apply (transfer, metis Abs-atoms-inverse Rep-atoms atom-map-compl-pres atom-map-top-pres diff-eq double-compl inf-le1 rangeE rangeI)

apply (transfer, smt (verit, ccfv-threshold) Abs-atoms-inverse Rep-atoms atom-map-inf-pres atom-map-sup-pres image-iff rangeI sup-inf-distrib1)

apply (transfer, metis (no-types, opaque-lifting) Abs-atoms-inverse Diff-disjoint Rep-atoms atom-map-compl-pres rangeE rangeI)

apply (transfer, smt Abs-atoms-inverse uminus-atoms.abs-eq Rep-atoms Un-Diff-cancel atom-map-compl-pres atom-map-inf-pres atom-map-minus-pres atom-map-sup-pres atom-map-top-pres diff-eq double-compl inf-compl-bot-right rangeE rangeI sup-commute sup-compl-top)

apply (transfer, smt Abs-atoms-inverse Rep-atoms atom-map-compl-pres atom-map-inf-pres

atom-map-minus-pres diff-eq rangeE rangeI
done

end

The homomorphism *atom-map* can then be restricted in its output type to the powerset boolean algebra.

lemma *at-map-bot-pres: at-map $\perp = \perp$*
by (*simp add: at-map-def atom-map-bot-pres bot-atoms.transfer*)

lemma *at-map-top-pres: at-map $\top = \top$*
by (*simp add: at-map-def atom-map-top-pres top-atoms.transfer*)

lemma *at-map-compl-pres: at-map \circ uminus = uminus \circ at-map*
unfolding *fun-eq-iff* **by** (*simp add: Abs-atoms-inverse at-map-def atom-map-compl-pres uminus-atoms.abs-eq*)

lemma *at-map-sup-pres: at-map $(x \sqcup y) = at-map x \sqcup at-map y$*
unfolding *at-map-def comp-def* **by** (*metis (mono-tags, lifting) Abs-atoms-inverse atom-map-sup-pres rangeI sup-atoms.transfer*)

lemma *at-map-inf-pres: at-map $(x \sqcap y) = at-map x \sqcap at-map y$*
unfolding *at-map-def comp-def* **by** (*metis (mono-tags, lifting) Abs-atoms-inverse atom-map-inf-pres inf-atoms.transfer rangeI*)

lemma *at-map-minus-pres: at-map $(x - y) = at-map x - at-map y$*
unfolding *at-map-def comp-def* **by** (*simp add: Abs-atoms-inverse atom-map-minus-pres minus-atoms.abs-eq*)

context *atomic-boolean-algebra*
begin

In atomic boolean algebras, *atom-map* is an embedding that maps atoms of the boolean algebra to those of the powerset boolean algebra. Analogous properties hold for *at-map*.

lemma *inj-atom-map: inj atom-map*

proof –

{fix *x y :: 'a*
assume *x \neq y*
hence *x \sqcap -y \neq $\perp \vee$ -x \sqcap y \neq \perp*
by (*auto simp: inf-shunt*)
hence $\exists z. atom z \wedge (z \leq x \sqcap -y \vee z \leq -x \sqcap y)$
using *atomicity* **by** *blast*
hence $\exists z. atom z \wedge ((z \in atom-map x \wedge \neg(z \in atom-map y)) \vee (\neg(z \in atom-map x) \wedge z \in atom-map y))$
unfolding *atom-def atom-map-def* **by** (*clarsimp, metis diff-eq inf.orderE diff-shunt-var*)
hence *atom-map x \neq atom-map y*
by *blast*}

```

thus ?thesis
  by (meson injI)
qed

```

```

lemma atom-map-atom-pres: atom x  $\implies$  atom-map x = {x}
  unfolding atom-def atom-map-def by (auto simp: bot-less dual-order.order-iff-strict)

```

```

lemma atom-map-atom-pres2: atom x  $\implies$  atom (atom-map x)
proof -

```

```

  assume atom x
  hence atom-map x = {x}
    by (simp add: atom-map-atom-pres)
  thus atom (atom-map x)
    using bounded-lattice-class.atom-def by auto
qed

```

end

```

lemma inj-at-map: inj (at-map::'a::atomic-boolean-algebra  $\implies$  'a atoms)
  unfolding at-map-def comp-def by (metis (no-types, lifting) Abs-atoms-inverse
inj-atom-map inj-def rangeI)

```

```

lemma at-map-atom-pres: atom (x::'a::atomic-boolean-algebra)  $\implies$  at-map x =
Abs-atoms {x}
  unfolding at-map-def comp-def by (simp add: atom-map-atom-pres)

```

```

lemma at-map-atom-pres2: atom (x::'a::atomic-boolean-algebra)  $\implies$  atom (at-map
x)
  unfolding at-map-def comp-def
  by (metis Abs-atoms-inverse atom-def atom-map-atom-pres2 atom-map-bot-pres
bot-atoms.abs-eq less-atoms.abs-eq rangeI)

```

Homomorphic images of atomic boolean algebras under atom-map are therefore atomic (rather obviously).

```

instance atoms :: (atomic-boolean-algebra) atomic-boolean-algebra

```

```

proof intro-classes

```

```

  fix x::'a atoms
  assume x  $\neq$   $\perp$ 
  hence  $\exists y. x = \text{at-map } y \wedge x \neq \perp$ 
    unfolding at-map-def comp-def by (metis Abs-atoms-cases rangeE)
  hence  $\exists y. x = \text{at-map } y \wedge (\exists z. \text{atom } z \wedge z \leq y)$ 
    using at-map-bot-pres atomicity by force
  hence  $\exists y. x = \text{at-map } y \wedge (\exists z. \text{atom } (\text{at-map } z) \wedge \text{at-map } z \leq \text{at-map } y)$ 
    by (metis at-map-atom-pres2 at-map-sup-pres sup.orderE sup-ge2)
  thus  $\exists y. \text{atom } y \wedge y \leq x$ 
    by fastforce
qed

```

```

context complete-boolean-algebra-alt

```

begin

In complete boolean algebras, atom-map is surjective; more precisely it is the left inverse of Sup, at least for sets of atoms. Below, this statement is made more explicit for at-map.

lemma *surj-atom-map*: $Y \subseteq \text{Collect atom} \implies Y = \text{atom-map} (\bigsqcup Y)$

proof

assume $Y \subseteq \text{Collect atom}$

thus $Y \subseteq \text{atom-map} (\bigsqcup Y)$

using *Sup-upper atom-map-def* **by** *force*

next

assume $Y \subseteq \text{Collect atom}$

hence $a: \forall y. y \in Y \longrightarrow \text{atom } y$

by *blast*

{fix z

assume $h: z \in \text{Collect atom} - Y$

hence $\forall y \in Y. y \sqcap z = \perp$

by (*metis DiffE a h atom-def dual-order.not-eq-order-implies-strict inf.absorb-iff2 inf-le2 inf-shunt mem-Collect-eq*)

hence $\bigsqcup Y \sqcap z = \perp$

using *Sup-least inf-shunt* **by** *simp*

hence $z \notin \text{atom-map} (\bigsqcup Y)$

using *atom-map-bot-pres atom-map-def* **by** *force*}

thus $\text{atom-map} (\bigsqcup Y) \subseteq Y$

using *atom-map-def* **by** *force*

qed

In this setting, atom-map is a complete boolean algebra morphism.

lemma *atom-map-Sup-pres*: $\text{atom-map} (\bigsqcup X) = (\bigcup x \in X. \text{atom-map } x)$

proof –

{fix z

have $(z \in \text{atom-map} (\bigsqcup X)) = (\text{atom } z \wedge z \leq \bigsqcup X)$

by (*simp add: atom-map-def*)

also have $\dots = (\text{atom } z \wedge (\exists x \in X. z \leq x))$

using *atom-Sup-iff* **by** *auto*

also have $\dots = (z \in (\bigcup x \in X. \text{atom-map } x))$

using *atom-map-def* **by** *auto*

finally have $(z \in \text{atom-map} (\bigsqcup X)) = (z \in (\bigcup x \in X. \text{atom-map } x))$

by *blast*}

thus *?thesis*

by *blast*

qed

lemma *atom-map-Sup-pres-var*: $\text{atom-map} \circ \text{Sup} = \text{Sup} \circ (\cdot) \text{atom-map}$

unfolding *fun-eq-iff comp-def* **by** (*simp add: atom-map-Sup-pres*)

For Inf-preservation, it is important that Infs are restricted to homomorphic images; hence they need to be pushed into the set of all atoms.

lemma *atom-map-Inf-pres*: $atom\text{-}map (\prod X) = Collect\ atom \cap (\bigcap x \in X. atom\text{-}map\ x)$

proof –

have $atom\text{-}map (\prod X) = atom\text{-}map (\neg(\bigsqcup x \in X. \neg x))$
by (*smt Collect-cong SUP-le-iff atom-map-def compl-le-compl-iff compl-le-swap1 le-Inf-iff*)
also have $\dots = Collect\ atom - atom\text{-}map (\bigsqcup x \in X. \neg x)$
using *atom-map-compl-pres* **by** *blast*
also have $\dots = Collect\ atom - (\bigcup x \in X. atom\text{-}map (\neg x))$
by (*simp add: atom-map-Sup-pres*)
also have $\dots = Collect\ atom - (\bigcup x \in X. Collect\ atom - atom\text{-}map\ x)$
using *atom-map-compl-pres* **by** *blast*
also have $\dots = Collect\ atom \cap (\bigcap x \in X. atom\text{-}map\ x)$
by *blast*
finally show *?thesis*.
qed

end

It follows that homomorphic images of complete boolean algebras under *atom-map* form complete boolean algebras.

instantiation *atoms* :: (*complete-boolean-algebra-alt*) *complete-boolean-algebra-alt atoms*
begin

lift-definition *Inf-atoms* :: '*a*::*complete-boolean-algebra-alt atoms set* \Rightarrow '*a*::*complete-boolean-algebra-alt atoms* is $\lambda X. Abs\text{-}atoms (Collect\ atom \cap Inter ((\cdot) Rep\text{-}atoms\ X))$.

lift-definition *Sup-atoms* :: '*a*::*complete-boolean-algebra-alt atoms set* \Rightarrow '*a*::*complete-boolean-algebra-alt atoms* is $\lambda X. Abs\text{-}atoms (Union ((\cdot) Rep\text{-}atoms\ X))$.

instance

apply (*intro-classes; transfer*)
apply (*metis (no-types, opaque-lifting) Abs-atoms-inverse image-iff inf-le1 le-Inf-iff le-infI2 order-refl rangeI surj-atom-map*)
apply (*metis (no-types, lifting) Abs-atoms-inverse Int-subset-iff Rep-atoms Sup-upper atom-map-atoms inf-le1 le-INF-iff rangeI surj-atom-map*)
apply (*metis Abs-atoms-inverse Rep-atoms SUP-least SUP-upper Sup-upper atom-map-atoms rangeI surj-atom-map*)
apply (*metis Abs-atoms-inverse Rep-atoms SUP-least Sup-upper atom-map-atoms rangeI surj-atom-map*)
by *simp-all*

end

Once more, properties proved above can now be restricted to *at-map*.

lemma *surj-at-map-var*: $at\text{-}map \circ Sup \circ Rep\text{-}atoms = (id::'a::complete-boolean-algebra-alt atoms \Rightarrow 'a\ atoms)$

unfolding *at-map-def comp-def fun-eq-iff id-def* **by** (*metis Rep-atoms Rep-atoms-inverse Sup-upper atom-map-atoms surj-atom-map*)

lemma *surj-at-map*: *surj* (*at-map*::'*a*::*complete-boolean-algebra-alt* \Rightarrow '*a atoms*)
unfolding *surj-def at-map-def comp-def* **by** (*metis Rep-atoms Rep-atoms-inverse image-iff*)

lemma *at-map-Sup-pres*: *at-map* \circ *Sup* = *Sup* \circ (\cdot) (*at-map*::'*a*::*complete-boolean-algebra-alt* \Rightarrow '*a atoms*)
unfolding *fun-eq-iff at-map-def comp-def atom-map-Sup-pres* **by** (*smt Abs-atoms-inverse Sup.SUP-cong Sup-atoms.transfer UN-extend-simps(10) rangeI*)

lemma *at-map-Sup-pres-var*: *at-map* (\bigsqcup *X*) = (\bigsqcup (*x*::'*a*::*complete-boolean-algebra-alt*) \in *X*. (*at-map x*))
using *at-map-Sup-pres comp-eq-elim* **by** *blast*

lemma *at-map-Inf-pres*: *at-map* (\prod *X*) = *Abs-atoms* (*Collect atom* \sqcap (\prod *x* \in *X*. (*Rep-atoms* (*at-map* (*x*::'*a*::*complete-boolean-algebra-alt*))))))
unfolding *at-map-def comp-def* **by** (*metis (no-types, lifting) Abs-atoms-inverse Sup.SUP-cong atom-map-Inf-pres rangeI*)

lemma *at-map-Inf-pres-var*: *at-map* \circ *Inf* = *Inf* \circ (\cdot) (*at-map*::'*a*::*complete-boolean-algebra-alt* \Rightarrow '*a atoms*)
unfolding *fun-eq-iff comp-def* **by** (*metis Inf-atoms.abs-eq at-map-Inf-pres image-image*)

Finally, on complete atomic boolean algebras (CABAs), *at-map* is an isomorphism, that is, a bijection that preserves the complete boolean algebra operations. Thus every CABA is isomorphic to a powerset boolean algebra and every powerset boolean algebra is a CABA. The bijective pair is given by *at-map* and *Sup* (defined on the powerset algebra). This theorem is a little version of Stone's theorem. In the general case, ultrafilters play the role of atoms.

lemma *Sup* \circ *atom-map* = (*id*::'*a*::*complete-atomic-boolean-algebra* \Rightarrow '*a*)
unfolding *fun-eq-iff comp-def id-def*
by (*metis Union-upper atom-map-atoms inj-atom-map inj-def rangeI surj-atom-map*)

lemma *inj-at-map-var*: *Sup* \circ *Rep-atoms* \circ *at-map* = (*id*::'*a*::*complete-atomic-boolean-algebra* \Rightarrow '*a*)

unfolding *at-map-def comp-def fun-eq-iff id-def*
by (*metis Abs-atoms-inverse Union-upper atom-map-atoms inj-atom-map inj-def rangeI surj-atom-map*)

lemma *bij-at-map*: *bij* (*at-map*::'*a*::*complete-atomic-boolean-algebra* \Rightarrow '*a atoms*)
unfolding *bij-def* **by** (*simp add: inj-at-map surj-at-map*)

instance *atoms* :: (*complete-atomic-boolean-algebra*) *complete-atomic-boolean-algebra..*

A full consideration of Stone duality is left for future work.

end

6 Galois Connections

theory *Galois-Connections*
imports *Order-Lattice-Props*

begin

6.1 Definitions and Basic Properties

The approach follows the Compendium of Continuous Lattices [3], without attempting completeness. First, left and right adjoints of a Galois connection are defined.

definition $adj :: ('a::ord \Rightarrow 'b::ord) \Rightarrow ('b \Rightarrow 'a) \Rightarrow bool$ (**infixl** \dashv 70) **where**
 $(f \dashv g) = (\forall x y. (f x \leq y) = (x \leq g y))$

definition $ladj (g::'a::Inf \Rightarrow 'b::ord) = (\lambda x. \sqcap \{y. x \leq g y\})$

definition $radj (f::'a::Sup \Rightarrow 'b::ord) = (\lambda y. \sqcup \{x. f x \leq y\})$

lemma *ladj-radj-dual*:

fixes $f :: 'a::complete-lattice-with-dual \Rightarrow 'b::ord-with-dual$

shows $ladj f x = \partial (radj (\partial_F f) (\partial x))$

proof –

have $ladj f x = \partial (\sqcup \{\partial y. \partial (f y) \leq \partial x\})$

unfolding *ladj-def* **by** (*metis* (*no-types*, *lifting*) *Collect-cong Inf-dual-var dual-dual-ord dual-iff*)

also have $\dots = \partial (\sqcup \{\partial y | y. \partial (f y) \leq \partial x\})$

by (*simp add: setcompr-eq-image*)

ultimately show *?thesis*

unfolding *ladj-def radj-def map-dual-def comp-def*

by (*smt* (*verit*) *Collect-cong invol-dual-var*)

qed

lemma *radj-ladj-dual*:

fixes $f :: 'a::complete-lattice-with-dual \Rightarrow 'b::ord-with-dual$

shows $radj f x = \partial (ladj (\partial_F f) (\partial x))$

by (*metis fun-dual5 invol-dual-var ladj-radj-dual map-dual-def*)

lemma *ladj-prop*:

fixes $g :: 'b::Inf \Rightarrow 'a::ord-with-dual$

shows $ladj g = Inf \circ (-') g \circ \uparrow$

unfolding *ladj-def vimage-def upset-prop fun-eq-iff comp-def* **by** *simp*

lemma *radj-prop*:

fixes $f :: 'b::Sup \Rightarrow 'a::ord$

shows $radj f = Sup \circ (-') f \circ \downarrow$

unfolding *radj-def vimage-def downset-prop fun-eq-iff comp-def* **by** *simp*

The first set of properties holds without any sort assumptions.

lemma *adj-iso1*: $f \dashv g \implies \text{mono } f$
unfolding *adj-def mono-def* **by** (*meson dual-order.refl dual-order.trans*)

lemma *adj-iso2*: $f \dashv g \implies \text{mono } g$
unfolding *adj-def mono-def* **by** (*meson dual-order.refl dual-order.trans*)

lemma *adj-comp*: $f \dashv g \implies \text{adj } h \ k \implies (f \circ h) \dashv (k \circ g)$
by (*simp add: adj-def*)

lemma *adj-dual*:
fixes $f :: 'a::\text{ord-with-dual} \Rightarrow 'b::\text{ord-with-dual}$
shows $f \dashv g = (\partial_F g) \dashv (\partial_F f)$
unfolding *adj-def map-dual-def comp-def* **by** (*metis (mono-tags, opaque-lifting) dual-dual-ord invol-dual-var*)

6.2 Properties for (Pre)Orders

The next set of properties holds in preorders or orders.

lemma *adj-cancel1*:
fixes $f :: 'a::\text{preorder} \Rightarrow 'b::\text{ord}$
shows $f \dashv g \implies f \circ g \leq \text{id}$
by (*simp add: adj-def le-funI*)

lemma *adj-cancel2*:
fixes $f :: 'a::\text{ord} \Rightarrow 'b::\text{preorder}$
shows $f \dashv g \implies \text{id} \leq g \circ f$
by (*simp add: adj-def eq-iff le-funI*)

lemma *adj-prop*:
fixes $f :: 'a::\text{preorder} \Rightarrow 'a$
shows $f \dashv g \implies f \circ g \leq g \circ f$
using *adj-cancel1 adj-cancel2 order-trans* **by** *blast*

lemma *adj-cancel-eq1*:
fixes $f :: 'a::\text{preorder} \Rightarrow 'b::\text{order}$
shows $f \dashv g \implies f \circ g \circ f = f$
unfolding *adj-def comp-def fun-eq-iff* **by** (*meson eq-iff order-refl order-trans*)

lemma *adj-cancel-eq2*:
fixes $f :: 'a::\text{order} \Rightarrow 'b::\text{preorder}$
shows $f \dashv g \implies g \circ f \circ g = g$
unfolding *adj-def comp-def fun-eq-iff* **by** (*meson eq-iff order-refl order-trans*)

lemma *adj-idem1*:
fixes $f :: 'a::\text{preorder} \Rightarrow 'b::\text{order}$
shows $f \dashv g \implies (f \circ g) \circ (f \circ g) = f \circ g$

by (*simp add: adj-cancel-eq1 rewriteL-comp-comp*)

lemma *adj-idem2*:

fixes $f :: 'a::order \Rightarrow 'b::preorder$

shows $f \dashv g \Longrightarrow (g \circ f) \circ (g \circ f) = g \circ f$

by (*simp add: adj-cancel-eq2 rewriteL-comp-comp*)

lemma *adj-iso3*:

fixes $f :: 'a::order \Rightarrow 'b::order$

shows $f \dashv g \Longrightarrow \text{mono } (f \circ g)$

by (*simp add: adj-iso1 adj-iso2 monoD monoI*)

lemma *adj-iso4*:

fixes $f :: 'a::order \Rightarrow 'b::order$

shows $f \dashv g \Longrightarrow \text{mono } (g \circ f)$

by (*simp add: adj-iso1 adj-iso2 monoD monoI*)

lemma *adj-canc1*:

fixes $f :: 'a::order \Rightarrow 'b::ord$

shows $f \dashv g \Longrightarrow ((f \circ g) x = (f \circ g) y \longrightarrow g x = g y)$

unfolding *adj-def comp-def* **by** (*metis eq-iff*)

lemma *adj-canc2*:

fixes $f :: 'a::ord \Rightarrow 'b::order$

shows $f \dashv g \Longrightarrow ((g \circ f) x = (g \circ f) y \longrightarrow f x = f y)$

unfolding *adj-def comp-def* **by** (*metis eq-iff*)

lemma *adj-sur-inv*:

fixes $f :: 'a::preorder \Rightarrow 'b::order$

shows $f \dashv g \Longrightarrow ((\text{surj } f) = (f \circ g = \text{id}))$

unfolding *adj-def surj-def comp-def* **by** (*metis eq-id-iff eq-iff order-refl order-trans*)

lemma *adj-surj-inj*:

fixes $f :: 'a::order \Rightarrow 'b::order$

shows $f \dashv g \Longrightarrow ((\text{surj } f) = (\text{inj } g))$

unfolding *adj-def inj-def surj-def* **by** (*metis eq-iff order-trans*)

lemma *adj-inj-inv*:

fixes $f :: 'a::preorder \Rightarrow 'b::order$

shows $f \dashv g \Longrightarrow ((\text{inj } f) = (g \circ f = \text{id}))$

by (*metis adj-cancel-eq1 eq-id-iff inj-def o-apply*)

lemma *adj-inj-surj*:

fixes $f :: 'a::order \Rightarrow 'b::order$

shows $f \dashv g \Longrightarrow ((\text{inj } f) = (\text{surj } g))$

unfolding *adj-def inj-def surj-def* **by** (*metis eq-iff order-trans*)

lemma *surj-id-the-inv*: $\text{surj } f \Longrightarrow g \circ f = \text{id} \Longrightarrow g = \text{the-inv } f$

by (*metis comp-apply id-apply inj-on-id inj-on-imageI2 surj-fun-eq the-inv-f-f*)


```

lemma inj-id-the-inv:  $\text{inj } f \implies f \circ g = \text{id} \implies f = \text{the-inv } g$ 
proof –
  assume a1:  $\text{inj } f$ 
  assume  $f \circ g = \text{id}$ 
  hence  $\forall x. \text{the-inv } g \ x = f \ x$ 
  using a1 by (metis (no-types) comp-apply eq-id-iff inj-on-id inj-on-imageI2
the-inv-f-f)
  thus ?thesis
  by presburger
qed

```

6.3 Properties for Complete Lattices

The next laws state that a function between complete lattices preserves infs if and only if it has a lower adjoint.

```

lemma radj-Inf-pres:
  fixes  $g :: 'b::\text{complete-lattice} \Rightarrow 'a::\text{complete-lattice}$ 
  shows  $(\exists f. f \dashv g) \implies \text{Inf-pres } g$ 
  apply (rule antisym, simp-all add: le-fun-def adj-def, safe)
  apply (meson INF-greatest Inf-lower dual-order.refl dual-order.trans)
  by (meson Inf-greatest dual-order.refl le-INF-iff)

```

```

lemma ladj-Sup-pres:
  fixes  $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'b::\text{complete-lattice-with-dual}$ 
  shows  $(\exists g. f \dashv g) \implies \text{Sup-pres } f$ 
  using Sup-pres-map-dual-var adj-dual radj-Inf-pres by blast

```

```

lemma radj-adj:
  fixes  $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$ 
  shows  $f \dashv g \implies g = (\text{radj } f)$ 
  unfolding adj-def radj-def by (metis (mono-tags, lifting) cSup-eq-maximum eq-iff
mem-Collect-eq)

```

```

lemma ladj-adj:
  fixes  $g :: 'b::\text{complete-lattice-with-dual} \Rightarrow 'a::\text{complete-lattice-with-dual}$ 
  shows  $f \dashv g \implies f = (\text{ladj } g)$ 
  unfolding adj-def ladj-def by (metis (no-types, lifting) cInf-eq-minimum eq-iff
mem-Collect-eq)

```

```

lemma Inf-pres-radj-aux:
  fixes  $g :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$ 
  shows  $\text{Inf-pres } g \implies (\text{ladj } g) \dashv g$ 
proof –
  assume a:  $\text{Inf-pres } g$ 
  {fix  $x \ y$ 
  assume b:  $\text{ladj } g \ x \leq y$ 
  hence  $g (\text{ladj } g \ x) \leq g \ y$ 
  by (simp add: Inf-subdistl-iso a monoD)

```

hence $\sqcap\{g\ y \mid y. x \leq g\ y\} \leq g\ y$
by (*metis a comp-eq-dest-lhs setcompr-eq-image ladj-def*)
hence $x \leq g\ y$
using *dual-order.trans le-Inf-iff* **by** *blast*
hence $ladj\ g\ x \leq y \longrightarrow x \leq g\ y$
by *simp*
thus *?thesis*
unfolding *adj-def ladj-def* **by** (*meson CollectI Inf-lower*)
qed

lemma *Sup-pres-ladj-aux*:
fixes $f :: 'a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual$
shows $Sup-pres\ f \Longrightarrow f \dashv (radj\ f)$
by (*metis (no-types, opaque-lifting) Inf-pres-radj-aux Sup-pres-map-dual-var adj-dual fun-dual5 map-dual-def radj-adj*)

lemma *Inf-pres-radj*:
fixes $g :: 'b::complete-lattice \Rightarrow 'a::complete-lattice$
shows $Inf-pres\ g \Longrightarrow (\exists f. f \dashv g)$
using *Inf-pres-radj-aux* **by** *fastforce*

lemma *Sup-pres-ladj*:
fixes $f :: 'a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual$
shows $Sup-pres\ f \Longrightarrow (\exists g. f \dashv g)$
using *Sup-pres-ladj-aux* **by** *fastforce*

lemma *Inf-pres-upper-adj-eq*:
fixes $g :: 'b::complete-lattice \Rightarrow 'a::complete-lattice$
shows $(Inf-pres\ g) = (\exists f. f \dashv g)$
using *radj-Inf-pres Inf-pres-radj* **by** *blast*

lemma *Sup-pres-ladj-eq*:
fixes $f :: 'a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual$
shows $(Sup-pres\ f) = (\exists g. f \dashv g)$
using *Sup-pres-ladj ladj-Sup-pres* **by** *blast*

lemma *Sup-downset-adj*: $(Sup::'a::complete-lattice\ set \Rightarrow 'a) \dashv \downarrow$
unfolding *adj-def downset-prop Sup-le-iff* **by** *force*

lemma *Sup-downset-adj-var*: $(Sup\ (X::'a::complete-lattice\ set) \leq y) = (X \subseteq \downarrow y)$
using *Sup-downset-adj adj-def* **by** *auto*

Once again many statements arise by duality, which Isabelle usually picks up.

end

7 Fixpoint Fusion

theory *Fixpoint-Fusion*

imports *Galois-Connections*

begin

Least and greatest fixpoint fusion laws for adjoints in a Galois connection, including some variants, are proved in this section. Again, the laws for least and greatest fixpoints are duals.

lemma *lfp-Fix*:

fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$

shows $\text{mono } f \implies \text{lfp } f = \bigsqcap (\text{Fix } f)$

unfolding *lfp-def Fix-def*

apply (*rule antisym*)

apply (*simp add: Collect-mono Inf-superset-mono*)

by (*metis (mono-tags) Inf-lower lfp-def lfp-unfold mem-Collect-eq*)

lemma *gfp-Fix*:

fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$

shows $\text{mono } f \implies \text{gfp } f = \bigsqcup (\text{Fix } f)$

by (*simp add: iso-map-dual gfp-to-lfp lfp-Fix Fix-map-dual-var Sup-to-Inf-var*)

lemma *gfp-little-fusion*:

fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'a$

and $g :: 'b::\text{complete-lattice} \Rightarrow 'b$

assumes *mono f*

assumes $h \circ f \leq g \circ h$

shows $h (\text{gfp } f) \leq \text{gfp } g$

proof –

have $h (f (\text{gfp } f)) \leq g (h (\text{gfp } f))$

using *assms(2) le-funD* **by** *fastforce*

hence $h (\text{gfp } f) \leq g (h (\text{gfp } f))$

by (*simp add: assms(1) gfp-fixpoint*)

thus $h (\text{gfp } f) \leq \text{gfp } g$

by (*simp add: gfp-upperbound*)

qed

lemma *lfp-little-fusion*:

fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$

and $g :: 'b::\text{complete-lattice-with-dual} \Rightarrow 'b$

assumes *mono f*

assumes $g \circ h \leq h \circ f$

shows $\text{lfp } g \leq h (\text{lfp } f)$

proof –

have $a: \text{mono } (\text{map-dual } f)$

by (*simp add: assms iso-map-dual*)

have $\text{map-dual } h \circ \text{map-dual } f \leq \text{map-dual } g \circ \text{map-dual } h$

by (*metis assms map-dual-anti map-dual-func1*)

thus *?thesis*

by (*metis a comp-eq-elim dual-dual-ord fun-dual1 gfp-little-fusion lfp-dual-var map-dual-def*)

qed

lemma *gfp-fusion*:

fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'a$

and $g :: 'b::\text{complete-lattice} \Rightarrow 'b$

assumes $\exists f. f \dashv h$

and *mono* f

and *mono* g

and $h \circ f = g \circ h$

shows $h (gfp\ f) = gfp\ g$

proof –

have $a: h (gfp\ f) \leq gfp\ g$

by (*simp add: assms(2) assms(4) gfp-little-fusion*)

obtain *hl* **where** *conn*: $\forall x\ y. (hl\ x \leq y) \longleftrightarrow (x \leq h\ y)$

using *assms adj-def* **by** *blast*

have $hl \circ g \leq hl \circ g \circ h \circ hl$

by (*simp add: le-fun-def, meson conn assms(3) monoE order-refl order-trans*)

also have $\dots = hl \circ h \circ f \circ hl$

by (*simp add: assms(4) comp-assoc*)

finally have $hl \circ g \leq f \circ hl$

by (*simp add: le-fun-def,metis conn inf.coboundedI2 inf.orderE order-refl*)

hence $hl (gfp\ g) \leq f (hl (gfp\ g))$

by (*metis comp-eq-dest-lhs gfp-unfold assms(3) le-fun-def*)

hence $hl (gfp\ g) \leq gfp\ f$

by (*simp add: gfp-upperbound*)

hence $gfp\ g \leq h (gfp\ f)$

by (*simp add: conn*)

thus *?thesis*

by (*simp add: a eq-iff*)

qed

lemma *lfp-fusion*:

fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$

and $g :: 'b::\text{complete-lattice-with-dual} \Rightarrow 'b$

assumes $\exists f. h \dashv f$

and *mono* f

and *mono* g

and $h \circ f = g \circ h$

shows $h (lfp\ f) = lfp\ g$

proof –

have $a: \exists f. \text{map-dual}\ f \dashv \text{map-dual}\ h$

using *adj-dual assms(1)* **by** *auto*

have $b: \text{mono} (\text{map-dual}\ f)$

by (*simp add: assms(2) iso-map-dual*)

have $c: \text{mono} (\text{map-dual}\ g)$

by (*simp add: assms(3) iso-map-dual*)

have $\text{map-dual}\ h \circ \text{map-dual}\ f = \text{map-dual}\ g \circ \text{map-dual}\ h$

by (*metis assms(4) map-dual-func1*)

thus *?thesis*

by (*metis a adj-dual b c gfp-fusion ladj-adj ladj-radj-dual lfp-dual-var lfp-to-gfp-var radj-adj*)
qed

lemma *gfp-fusion-inf-pres*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'a$
and $g :: 'b::\text{complete-lattice} \Rightarrow 'b$
assumes *Inf-pres h*
and *mono f*
and *mono g*
and $h \circ f = g \circ h$
shows $h (\text{gfp } f) = \text{gfp } g$
by (*simp add: Inf-pres-radj assms gfp-fusion*)

lemma *lfp-fusion-sup-pres*:
fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$
and $g :: 'b::\text{complete-lattice-with-dual} \Rightarrow 'b$
assumes *Sup-pres h*
and *mono f*
and *mono g*
and $h \circ f = g \circ h$
shows $h (\text{lfp } f) = \text{lfp } g$
by (*simp add: Sup-pres-ladj assms lfp-fusion*)

The following facts are useful for the semantics of isotone predicate transformers. A dual statement for least fixpoints can be proved, but is not spelled out here.

lemma *k-adju*:
fixes $k :: 'a::\text{order} \Rightarrow 'b::\text{complete-lattice}$
shows $\exists F. \forall x. (F :: 'b \Rightarrow 'a \Rightarrow 'b) \dashv (\lambda k. k y)$
by (*force intro!: fun-eq-iff Inf-pres-radj*)

lemma *k-adju-var*: $\exists F. \forall x. \forall f :: 'a::\text{order} \Rightarrow 'b::\text{complete-lattice}. (F x \leq f) = (x \leq (\lambda k. k y) f)$
using *k-adju unfolding adj-def by simp*

lemma *gfp-fusion-var*:
fixes $F :: ('a::\text{order} \Rightarrow 'b::\text{complete-lattice}) \Rightarrow 'a \Rightarrow 'b$
and $g :: 'b \Rightarrow 'b$
assumes *mono F*
and *mono g*
and $\forall h. F h x = g (h x)$
shows $\text{gfp } F x = \text{gfp } g$
by (*metis (no-types, opaque-lifting) assms eq-iff gfp-fixpoint gfp-upperbound k-adju-var monoE order-refl*)

This time, Isabelle is picking up dualities rather inconsistently.
end

8 Closure and Co-Closure Operators

```
theory Closure-Operators
  imports Galois-Connections
```

```
begin
```

8.1 Closure Operators

Closure and coclosure operators in orders and complete lattices are defined in this section, and some basic properties are proved. Isabelle infers the appropriate types. Facts are taken mainly from the Compendium of Continuous Lattices [3] and Rosenthal's book on quantales [10].

```
definition clop :: ('a::order  $\Rightarrow$  'a)  $\Rightarrow$  bool where
  clop f = (id  $\leq$  f  $\wedge$  mono f  $\wedge$  f  $\circ$  f  $\leq$  f)
```

```
lemma clop-extensive: clop f  $\Longrightarrow$  id  $\leq$  f
by (simp add: clop-def)
```

```
lemma clop-extensive-var: clop f  $\Longrightarrow$  x  $\leq$  f x
by (simp add: clop-def le-fun-def)
```

```
lemma clop-iso: clop f  $\Longrightarrow$  mono f
by (simp add: clop-def)
```

```
lemma clop-iso-var: clop f  $\Longrightarrow$  x  $\leq$  y  $\Longrightarrow$  f x  $\leq$  f y
by (simp add: clop-def mono-def)
```

```
lemma clop-idem: clop f  $\Longrightarrow$  f  $\circ$  f = f
by (simp add: antisym clop-def le-fun-def)
```

```
lemma clop-Fix-range: clop f  $\Longrightarrow$  (Fix f = range f)
by (simp add: clop-idem retraction-prop-fix)
```

```
lemma clop-idem-var: clop f  $\Longrightarrow$  f (f x) = f x
by (simp add: clop-idem retraction-prop)
```

```
lemma clop-Inf-closed-var:
fixes f :: 'a::complete-lattice  $\Rightarrow$  'a
shows clop f  $\Longrightarrow$  f  $\circ$  Inf  $\circ$  ( $\cdot$ ) f = Inf  $\circ$  ( $\cdot$ ) f
unfolding clop-def mono-def comp-def le-fun-def
by (metis (mono-tags, lifting) antisym id-apply le-INF-iff order-refl)
```

```
lemma clop-top:
fixes f :: 'a::complete-lattice  $\Rightarrow$  'a
shows clop f  $\Longrightarrow$  f  $\top$  =  $\top$ 
by (simp add: clop-extensive-var top.extremum-uniqueI)
```

```
lemma clop (f::'a::complete-lattice  $\Rightarrow$  'a)  $\Longrightarrow$  f ( $\bigsqcup$  x  $\in$  X. f x) = ( $\bigsqcup$  x  $\in$  X. f x)
```

oops

lemma *clop* ($f :: 'a :: \text{complete-lattice} \Rightarrow 'a$) $\implies f (f x \sqcup f y) = f x \sqcup f y$
oops

lemma *clop* ($f :: 'a :: \text{complete-lattice} \Rightarrow 'a$) $\implies f \perp = \perp$
oops

lemma *clop* ($f :: 'a \text{ set} \Rightarrow 'a \text{ set}$) $\implies f (\bigsqcup x \in X. f x) = (\bigsqcup x \in X. f x)$
oops

lemma *clop* ($f :: 'a \text{ set} \Rightarrow 'a \text{ set}$) $\implies f (f x \sqcup f y) = f x \sqcup f y$
oops

lemma *clop* ($f :: 'a \text{ set} \Rightarrow 'a \text{ set}$) $\implies f \perp = \perp$
oops

lemma *clop-closure*: $\text{clop } f \implies (x \in \text{range } f) = (f x = x)$
by (*simp add: clop-idem retraction-prop*)

lemma *clop-closure-set*: $\text{clop } f \implies \text{range } f = \text{Fix } f$
by (*simp add: clop-Fix-range*)

lemma *clop-closure-prop*: ($\text{clop} :: ('a :: \text{complete-lattice-with-dual} \Rightarrow 'a) \Rightarrow \text{bool}$) (*Inf*
 $\circ \uparrow$)
by (*simp add: clop-def mono-def*)

lemma *clop-closure-prop-var*: $\text{clop } (\lambda x :: 'a :: \text{complete-lattice}. \bigsqcap \{y. x \leq y\})$
unfolding *clop-def comp-def le-fun-def mono-def* **by** (*simp add: Inf-lower le-Inf-iff*)

lemma *clop-alt*: $(\text{clop } f) = (\forall x y. x \leq f y \longleftrightarrow f x \leq f y)$
unfolding *clop-def mono-def le-fun-def comp-def id-def* **by** (*meson dual-order.refl order-trans*)

Finally it is shown that adjoints in a Galois connection yield closure operators.

lemma *clop-adj*:
fixes $f :: 'a :: \text{order} \Rightarrow 'b :: \text{order}$
shows $f \dashv g \implies \text{clop } (g \circ f)$
by (*simp add: adj-cancel2 adj-idem2 adj-iso4 clop-def*)

Closure operators are monads for posets, and monads arise from adjunctions. This fact is not formalised at this point. But here is the first step: every function can be decomposed into a surjection followed by an injection.

definition *surj-on* $f Y = (\forall y \in Y. \exists x. y = f x)$

lemma *surj-surj-on*: $\text{surj } f \implies \text{surj-on } f Y$
by (*simp add: surjD surj-on-def*)

lemma *fun-surj-inj*: $\exists g h. f = g \circ h \wedge \text{surj-on } h (\text{range } f) \wedge \text{inj-on } g (\text{range } f)$
proof –
obtain *h* **where** *a*: $\forall x. f x = h x$
by *blast*
then have *surj-on h (range f)*
by (*metis (mono-tags, lifting) imageE surj-on-def*)
then show *?thesis*
unfolding *inj-on-def surj-on-def fun-eq-iff* **using** *a* **by** *auto*
qed

Connections between downsets, upsets and closure operators are outlined next.

lemma *preorder-clop*: $\text{clop } (\Downarrow :: 'a :: \text{preorder set} \Rightarrow 'a \text{ set})$
by (*simp add: clop-def downset-set-ext downset-set-iso*)

lemma *clop-preorder-aux*: $\text{clop } f \Longrightarrow (x \in f \{y\} \longleftrightarrow f \{x\} \subseteq f \{y\})$
by (*simp add: clop-alt*)

lemma *clop-preorder*: $\text{clop } f \Longrightarrow \text{class.preorder } (\lambda x y. f \{x\} \subseteq f \{y\}) (\lambda x y. f \{x\} \subset f \{y\})$
unfolding *clop-def mono-def le-fun-def id-def comp-def* **by** *standard (auto simp: subset-not-subset-eq)*

lemma *preorder-clop-dual*: $\text{clop } (\Uparrow :: 'a :: \text{preorder-with-dual set} \Rightarrow 'a \text{ set})$
by (*simp add: clop-def upset-set-anti upset-set-ext*)

The closed elements of any closure operator over a complete lattice form an Inf-closed set (a Moore family).

lemma *clop-Inf-closed*:
fixes *f* :: *'a*::*complete-lattice* \Rightarrow *'a*
shows $\text{clop } f \Longrightarrow \text{Inf-closed-set } (\text{Fix } f)$
unfolding *clop-def Inf-closed-set-def mono-def le-fun-def comp-def id-def Fix-def*
by (*smt (verit) Inf-greatest Inf-lower antisym mem-Collect-eq subsetCE*)

lemma *clop-top-Fix*:
fixes *f* :: *'a*::*complete-lattice* \Rightarrow *'a*
shows $\text{clop } f \Longrightarrow \top \in \text{Fix } f$
by (*simp add: clop-Fix-range clop-closure clop-top*)

Conversely, every Inf-closed subset of a complete lattice is the set of fixpoints of some closure operator.

lemma *Inf-closed-clop*:
fixes *X* :: *'a*::*complete-lattice set*
shows $\text{Inf-closed-set } X \Longrightarrow \text{clop } (\lambda y. \bigcap \{x \in X. y \leq x\})$
by (*smt (verit) Collect-mono-iff Inf-superset-mono clop-alt dual-order.trans le-Inf-iff mem-Collect-eq*)

lemma *Inf-closed-clop-var*:


```

fixes X :: 'a::complete-lattice set
shows clop f  $\implies \forall x \in X. x \in \text{range } f \implies \bigsqcap X \in \text{range } f$ 
by (metis Inf-closed-set-def clop-Fix-range clop-Inf-closed subsetI)

```

It is well known that downsets and upsets over an ordering form subalgebras of the complete powerset lattice.

```

typedef (overloaded) 'a downsets = range ( $\Downarrow$ ::'a::order set  $\Rightarrow$  'a set)
by fastforce

```

```

setup-lifting type-definition-downsets

```

```

typedef (overloaded) 'a upsets = range ( $\Uparrow$ ::'a::order set  $\Rightarrow$  'a set)
by fastforce

```

```

setup-lifting type-definition-upsets

```

```

instantiation downsets :: (order) Inf-lattice
begin

```

```

lift-definition Inf-downsets :: 'a downsets set  $\Rightarrow$  'a downsets is Abs-downsets  $\circ$ 
Inf  $\circ$  ( $\cdot$ ) Rep-downsets.

```

```

lift-definition less-eq-downsets :: 'a downsets  $\Rightarrow$  'a downsets  $\Rightarrow$  bool is  $\lambda X Y. \text{Rep-downsets } X \subseteq \text{Rep-downsets } Y$ .

```

```

lift-definition less-downsets :: 'a downsets  $\Rightarrow$  'a downsets  $\Rightarrow$  bool is  $\lambda X Y. \text{Rep-downsets } X \subset \text{Rep-downsets } Y$ .

```

```

instance

```

```

apply intro-classes
apply (transfer, simp)
apply (transfer, blast)
apply (simp add: Closure-Operators.less-eq-downsets.abs-eq Rep-downsets-inject)
apply (transfer, smt (verit) Abs-downsets-inverse INF-lower Inf-closed-clop-var
Rep-downsets image-iff o-def preorder-clop)
apply (transfer, smt (verit) comp-def Abs-downsets-inverse Inf-closed-clop-var
Rep-downsets image-iff le-INF-iff preorder-clop)
done

```

```

end

```

```

instantiation upsets :: (order-with-dual) Inf-lattice
begin

```

```

lift-definition Inf-upsets :: 'a upsets set  $\Rightarrow$  'a upsets is Abs-upsets  $\circ$  Inf  $\circ$  ( $\cdot$ )
Rep-upsets.

```

```

lift-definition less-eq-upsets :: 'a upsets  $\Rightarrow$  'a upsets  $\Rightarrow$  bool is  $\lambda X Y. \text{Rep-upsets } X \subseteq \text{Rep-upsets } Y$ .

```

lift-definition *less-upsets* :: 'a upsets \Rightarrow 'a upsets \Rightarrow bool **is** $\lambda X Y. \text{Rep-upsets } X \subset \text{Rep-upsets } Y$.

instance

```

apply intro-classes
  apply (transfer, simp)
  apply (transfer, blast)
  apply (simp add: Closure-Operators.less-eq-upsets.abs-eq Rep-upsets-inject)
  apply (transfer, smt (verit) Abs-upsets-inverse Inf-closed-clop-var Inf-lower
Rep-upsets comp-apply image-iff preorder-clop-dual)
  apply (transfer, smt (verit) comp-def Abs-upsets-inverse Inf-closed-clop-var In-
ter-iff Rep-upsets image-iff preorder-clop-dual subsetCE subsetI)
done

```

end

It has already been shown in the section on representations that the map ds , which maps elements of the order to its downset, is an order embedding. However, the duality between the underlying ordering and the lattices of up- and down-closed sets as categories can probably not be expressed, as there is no easy access to contravariant functors.

8.2 Co-Closure Operators

Next, the co-closure (or kernel) operation satisfies dual laws.

definition *coclop* :: ('a::order \Rightarrow 'a::order) \Rightarrow bool **where**
coclop $f = (f \leq id \wedge mono\ f \wedge f \leq f \circ f)$

lemma *coclop-dual*: (*coclop*::('a::order-with-dual \Rightarrow 'a) \Rightarrow bool) = *clop* $\circ \partial_F$
unfolding *coclop-def clop-def id-def mono-def map-dual-def comp-def fun-eq-iff*
le-fun-def
by (*metis invol-dual-var ord-dual*)

lemma *coclop-dual-var*:
fixes $f :: 'a::order-with-dual \Rightarrow 'a$
shows *coclop* $f = clop (\partial_F f)$
by (*simp add: coclop-dual*)

lemma *clop-dual*: (*clop*::('a::order-with-dual \Rightarrow 'a) \Rightarrow bool) = *coclop* $\circ \partial_F$
by (*simp add: coclop-dual comp-assoc map-dual-invol*)

lemma *clop-dual-var*:
fixes $f :: 'a::order-with-dual \Rightarrow 'a$
shows *clop* $f = coclop (\partial_F f)$
by (*simp add: clop-dual*)

lemma *coclop-coextensive*: *coclop* $f \Longrightarrow f \leq id$

by (*simp add: coclop-def*)

lemma *coclop-coextensive-var*: $\text{coclop } f \implies f x \leq x$
using *coclop-def le-funD* **by** *fastforce*

lemma *coclop-iso*: $\text{coclop } f \implies \text{mono } f$
by (*simp add: coclop-def*)

lemma *coclop-iso-var*: $\text{coclop } f \implies (x \leq y \longrightarrow f x \leq f y)$
by (*simp add: coclop-iso monoD*)

lemma *coclop-idem*: $\text{coclop } f \implies f \circ f = f$
by (*simp add: antisym coclop-def le-fun-def*)

lemma *coclop-closure*: $\text{coclop } f \implies (x \in \text{range } f) = (f x = x)$
by (*simp add: coclop-idem retraction-prop*)

lemma *coclop-Fix-range*: $\text{coclop } f \implies (\text{Fix } f = \text{range } f)$
by (*simp add: coclop-idem retraction-prop-fix*)

lemma *coclop-idem-var*: $\text{coclop } f \implies f (f x) = f x$
by (*simp add: coclop-idem retraction-prop*)

lemma *coclop-Sup-closed-var*:
fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$
shows $\text{coclop } f \implies f \circ \text{Sup} \circ (\cdot) f = \text{Sup} \circ (\cdot) f$
unfolding *coclop-def mono-def comp-def le-fun-def*
by (*metis (mono-tags, lifting) SUP-le-iff antisym id-apply order-refl*)

lemma *Sup-closed-coclop-var*:
fixes $X :: 'a::\text{complete-lattice set}$
shows $\text{coclop } f \implies \forall x \in X. x \in \text{range } f \implies \bigsqcup X \in \text{range } f$
by (*smt (verit) Inf.INF-id-eq Sup.SUP-cong antisym coclop-closure coclop-coextensive-var coclop-iso id-apply mono-SUP*)

lemma *coclop-bot*:
fixes $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'a$
shows $\text{coclop } f \implies f \perp = \perp$
by (*simp add: bot.extremum-uniqueI coclop-coextensive-var*)

lemma *coclop* ($f :: 'a::\text{complete-lattice} \Rightarrow 'a$) $\implies f (\bigsqcap x \in X. f x) = (\bigsqcap x \in X. f x)$
oops

lemma *coclop* ($f :: 'a::\text{complete-lattice} \Rightarrow 'a$) $\implies f (f x \sqcap f y) = f x \sqcap f y$
oops

lemma *coclop* ($f :: 'a::\text{complete-lattice} \Rightarrow 'a$) $\implies f \top = \top$
oops

lemma *coclop* ($f :: 'a \text{ set} \Rightarrow 'a \text{ set}$) $\Longrightarrow f (\bigsqcap x \in X. f x) = (\bigsqcap x \in X. f x)$
oops

lemma *coclop* ($f :: 'a \text{ set} \Rightarrow 'a \text{ set}$) $\Longrightarrow f (f x \sqcap f y) = f x \sqcap f y$
oops

lemma *coclop* ($f :: 'a \text{ set} \Rightarrow 'a \text{ set}$) $\Longrightarrow f \top = \top$
oops

lemma *coclop-coclosure*: *coclop* $f \Longrightarrow f x = x \longleftrightarrow x \in \text{range } f$
by (*simp add: coclop-idem retraction-prop*)

lemma *coclop-coclosure-set*: *coclop* $f \Longrightarrow \text{range } f = \text{Fix } f$
by (*simp add: coclop-idem retraction-prop-fix*)

lemma *coclop-coclosure-prop*: (*coclop*::('a::complete-lattice \Rightarrow 'a) \Rightarrow bool) (*Sup* \circ
 \downarrow)
by (*simp add: coclop-def mono-def*)

lemma *coclop-coclosure-prop-var*: *coclop* ($\lambda x :: 'a :: \text{complete-lattice}. \bigsqcup \{y. y \leq x\}$)
by (*metis (mono-tags, lifting) Sup-atMost atMost-def coclop-def comp-apply eq-id-iff eq-refl mono-def*)

lemma *coclop-alt*: (*coclop* f) = ($\forall x y. f x \leq y \longleftrightarrow f x \leq f y$)
unfolding *coclop-def mono-def le-fun-def comp-def id-def*
by (*meson dual-order.refl order-trans*)

lemma *coclop-adj*:
fixes $f :: 'a :: \text{order} \Rightarrow 'b :: \text{order}$
shows $f \dashv g \Longrightarrow \text{coclop } (f \circ g)$
by (*simp add: adj-cancel1 adj-idem1 adj-iso3 coclop-def*)

Finally, a subset of a complete lattice is Sup-closed if and only if it is the set of fixpoints of some co-closure operator.

lemma *coclop-Sup-closed*:
fixes $f :: 'a :: \text{complete-lattice} \Rightarrow 'a$
shows *coclop* $f \Longrightarrow \text{Sup-closed-set } (\text{Fix } f)$
unfolding *coclop-def Sup-closed-set-def mono-def le-fun-def comp-def id-def Fix-def*
by (*smt (verit) Sup-least Sup-upper antisym-conv mem-Collect-eq subsetCE*)

lemma *Sup-closed-coclop*:
fixes $X :: 'a :: \text{complete-lattice set}$
shows *Sup-closed-set* $X \Longrightarrow \text{coclop } (\lambda y. \bigsqcup \{x \in X. x \leq y\})$
unfolding *Sup-closed-set-def coclop-def mono-def le-fun-def comp-def*
apply *safe*
apply (*metis (no-types, lifting) Sup-least eq-id-iff mem-Collect-eq*)
apply (*smt (verit) Collect-mono-iff Sup-subset-mono dual-order.trans*)
by (*simp add: Collect-mono-iff Sup-subset-mono Sup-upper*)

8.3 Complete Lattices of Closed Elements

The machinery developed allows showing that the closed elements in a complete lattice (with respect to some closure operation) form themselves a complete lattice.

```

class cl-op = ord +
  fixes cl-op :: 'a ⇒ 'a
  assumes cl-op-ext: x ≤ cl-op x
  and cl-op-iso: x ≤ y ⇒ cl-op x ≤ cl-op y
  and cl-op-wtrans: cl-op (cl-op x) ≤ cl-op x

class clattice-with-clop = complete-lattice + cl-op

begin

lemma cl-op-cl-op: cl-op cl-op
  unfolding cl-op-def le-fun-def comp-def
  by (simp add: cl-op-class.cl-op-ext cl-op-class.cl-op-iso cl-op-class.cl-op-wtrans or-
der-class.mono-def)

lemma cl-op-idem [simp]: cl-op ∘ cl-op = cl-op
  using cl-op-ext cl-op-wtrans order.antisym by auto

lemma cl-op-idem-var [simp]: cl-op (cl-op x) = cl-op x
  by (simp add: order.antisym cl-op-ext cl-op-wtrans)

lemma cl-op-range-Fix: range cl-op = Fix cl-op
  by (simp add: retraction-prop-fix)

lemma Inf-closed-cl-op-var:
  fixes X :: 'a set
  shows ∀ x ∈ X. x ∈ range cl-op ⇒ ⋂ X ∈ range cl-op
proof -
  assume h: ∀ x ∈ X. x ∈ range cl-op
  hence ∀ x ∈ X. cl-op x = x
    by (simp add: retraction-prop)
  hence cl-op (⋂ X) = ⋂ X
    by (metis Inf-lower cl-op-ext cl-op-iso dual-order.antisym le-Inf-iff)
  thus ?thesis
    by (metis rangeI)
qed

lemma inf-closed-cl-op-var: x ∈ range cl-op ⇒ y ∈ range cl-op ⇒ x ⋂ y ∈ range
cl-op
  by (smt (verit) Inf-closed-cl-op-var UnI1 insert-iff insert-is-Un inf-Inf)

end

typedef (overloaded) 'a::clattice-with-clop cl-op-im = range (cl-op::'a ⇒ 'a)

```

by *force*

setup-lifting *type-definition-cl-op-im*

lemma *cl-op-prop [iff]*: $(cl\text{-}op\ (x \sqcup y) = cl\text{-}op\ y) = (cl\text{-}op\ (x :: 'a :: clattice\text{-}with\text{-}clop) \leq cl\text{-}op\ y)$
by (*smt (verit) cl-op-class.clop-iso clop-ext clop-wtrans inf-sup-ord(4) le-iff-sup sup.absorb-iff1 sup-left-commute*)

lemma *cl-op-prop-var [iff]*: $(cl\text{-}op\ (x \sqcup cl\text{-}op\ y) = cl\text{-}op\ y) = (cl\text{-}op\ (x :: 'a :: clattice\text{-}with\text{-}clop) \leq cl\text{-}op\ y)$
by (*metis cl-op-prop clattice-with-clop-class.clop-idem-var*)

instantiation *cl-op-im* :: $(clattice\text{-}with\text{-}clop)$ *complete-lattice*
begin

lift-definition *Inf-cl-op-im* :: $'a\ cl\text{-}op\text{-}im\ set \Rightarrow 'a\ cl\text{-}op\text{-}im$ **is** *Inf*
by (*simp add: Inf-closed-cl-op-var*)

lift-definition *Sup-cl-op-im* :: $'a\ cl\text{-}op\text{-}im\ set \Rightarrow 'a\ cl\text{-}op\text{-}im$ **is** $\lambda X. cl\text{-}op\ (\sqcup X)$
by *simp*

lift-definition *inf-cl-op-im* :: $'a\ cl\text{-}op\text{-}im \Rightarrow 'a\ cl\text{-}op\text{-}im \Rightarrow 'a\ cl\text{-}op\text{-}im$ **is** *inf*
by (*simp add: inf-closed-cl-op-var*)

lift-definition *sup-cl-op-im* :: $'a\ cl\text{-}op\text{-}im \Rightarrow 'a\ cl\text{-}op\text{-}im \Rightarrow 'a\ cl\text{-}op\text{-}im$ **is** $\lambda x\ y. cl\text{-}op\ (x \sqcup y)$
by *simp*

lift-definition *less-eq-cl-op-im* :: $'a\ cl\text{-}op\text{-}im \Rightarrow 'a\ cl\text{-}op\text{-}im \Rightarrow bool$ **is** (\leq) .

lift-definition *less-cl-op-im* :: $'a\ cl\text{-}op\text{-}im \Rightarrow 'a\ cl\text{-}op\text{-}im \Rightarrow bool$ **is** $(<)$.

lift-definition *bot-cl-op-im* :: $'a\ cl\text{-}op\text{-}im$ **is** $cl\text{-}op\ \perp$
by *simp*

lift-definition *top-cl-op-im* :: $'a\ cl\text{-}op\text{-}im$ **is** \top
by (*simp add: clop-cl-op clop-closure clop-top*)

instance
apply (*intro-classes; transfer*)
apply (*simp-all add: less-le-not-le Inf-lower Inf-greatest*)
apply (*meson clop-cl-op clop-extensive-var dual-order.trans inf-sup-ord(3)*)
apply (*meson clop-cl-op clop-extensive-var dual-order.trans sup-ge2*)
apply (*metis cl-op-class.clop-iso clop-cl-op clop-closure le-sup-iff*)
apply (*meson Sup-upper clop-cl-op clop-extensive-var dual-order.trans*)
by (*metis Sup-le-iff cl-op-class.clop-iso clop-cl-op clop-closure*)

end

This statement is perhaps less useful as it might seem, because it is difficult to make it cooperate with concrete closure operators, which one would not generally like to define within a type class. Alternatively, a sublocale statement could perhaps be given. It would also have been nice to prove this statement for Sup-lattices—this would have cut down the number of proof obligations significantly. But this would require a tighter integration of these structures. A similar statement could have been proved for co-closure operators. But this would not lead to new insights.

Next I show that for every surjective Sup-preserving function between complete lattices there is a closure operator such that the set of closed elements is isomorphic to the range of the surjection.

```
lemma surj-Sup-pres-id:  
  fixes  $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'b::\text{complete-lattice-with-dual}$   
  assumes surj f  
  and Sup-pres f  
  shows  $f \circ (\text{radj } f) = \text{id}$   
proof –  
  have  $f \dashv (\text{radj } f)$   
    using Sup-pres-ladj assms(2) radj-adj by auto  
  thus ?thesis  
    using adj-sur-inv assms(1) by blast  
qed
```

```
lemma surj-Sup-pres-inj:  
  fixes  $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'b::\text{complete-lattice-with-dual}$   
  assumes surj f  
  and Sup-pres f  
  shows  $\text{inj } (\text{radj } f)$   
  by (metis assms comp-eq-dest-lhs id-apply injI surj-Sup-pres-id)
```

```
lemma surj-Sup-pres-inj-on:  
  fixes  $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'b::\text{complete-lattice-with-dual}$   
  assumes surj f  
  and Sup-pres f  
  shows  $\text{inj-on } f (\text{range } (\text{radj } f \circ f))$   
  by (smt (verit) Sup-pres-ladj-aux adj-idem2 assms(2) comp-apply inj-on-def re-  
traction-prop)
```

```
lemma surj-Sup-pres-bij-on:  
  fixes  $f :: 'a::\text{complete-lattice-with-dual} \Rightarrow 'b::\text{complete-lattice-with-dual}$   
  assumes surj f  
  and Sup-pres f  
  shows  $\text{bij-betw } f (\text{range } (\text{radj } f \circ f)) \text{ UNIV}$   
  unfolding bij-betw-def  
  apply safe
```

```

apply (simp add: assms(1) assms(2) surj-Sup-pres-inj-on cong del: image-cong-simp)
apply auto
apply (metis (mono-tags) UNIV-I assms(1) assms(2) comp-apply id-apply im-
age-image surj-Sup-pres-id surj-def)
done

```

Thus the restriction of f to the set of closed elements is indeed a bijection. The final fact shows that it preserves Sups of closed elements, and hence is an isomorphism of complete lattices.

```

lemma surj-Sup-pres-iso:
  fixes  $f :: 'a::complete-lattice-with-dual \Rightarrow 'b::complete-lattice-with-dual$ 
  assumes surj  $f$ 
  and Sup-pres  $f$ 
  shows  $f ((\text{radj } f \circ f) (\bigsqcup X)) = (\bigsqcup x \in X. f x)$ 
  by (metis assms(1) assms(2) comp-def pointfree-idE surj-Sup-pres-id)

```

8.4 A Quick Example: Dedekind-MacNeille Completions

I only outline the basic construction. Additional facts about join density, and that the completion yields the least complete lattice that contains all Sups and Infs of the underlying posets, are left for future consideration.

abbreviation $dm \equiv lb\text{-set} \circ ub\text{-set}$

```

lemma up-set-prop:  $(X::'a::preorder\ set) \neq \{\}$   $\implies ub\text{-set } X = \bigcap \{\uparrow x \mid x. x \in X\}$ 
  unfolding ub-set-def upset-def upset-set-def by (safe, simp-all, blast)

```

```

lemma lb-set-prop:  $(X::'a::preorder\ set) \neq \{\}$   $\implies lb\text{-set } X = \bigcap \{\downarrow x \mid x. x \in X\}$ 
  unfolding lb-set-def downset-def downset-set-def by (safe, simp-all, blast)

```

```

lemma dm-downset-var:  $dm \{x\} = \downarrow(x::'a::preorder)$ 
  unfolding lb-set-def ub-set-def downset-def downset-set-def
  by (clarsimp, meson order-refl order-trans)

```

```

lemma dm-downset:  $dm \circ \eta = (\downarrow::'a::preorder \Rightarrow 'a\ set)$ 
  using dm-downset-var fun.map-cong by fastforce

```

```

lemma dm-inj:  $inj ((dm::'a::order\ set \Rightarrow 'a\ set) \circ \eta)$ 
  by (simp add: dm-downset downset-inj)

```

```

lemma clop (lb-set  $\circ$  ub-set)
  unfolding clop-def lb-set-def ub-set-def
  apply safe
  unfolding le-fun-def comp-def id-def mono-def
  by auto

```

end

9 Locale-Based Duality

```

theory Order-Lattice-Props-Loc
  imports Main
begin

```

```

unbundle lattice-syntax

```

This section explores order and lattice duality based on locales. Used within the context of a class or locale, this is very effective, though more opaque than the previous approach. Outside of such a context, however, it apparently cannot be used for dualising theorems. Examples are properties of functions between orderings or lattices.

```

definition Fix :: ('a ⇒ 'a) ⇒ 'a set where
  Fix f = {x. f x = x}

```

```

context ord
begin

```

```

definition min-set :: 'a set ⇒ 'a set where
  min-set X = {y ∈ X. ∀ x ∈ X. x ≤ y → x = y}

```

```

definition max-set :: 'a set ⇒ 'a set where
  max-set X = {x ∈ X. ∀ y ∈ X. x ≤ y → x = y}

```

```

definition directed :: 'a set ⇒ bool where
  directed X = (∀ Y. finite Y ∧ Y ⊆ X → (∃ x ∈ X. ∀ y ∈ Y. y ≤ x))

```

```

definition filtered :: 'a set ⇒ bool where
  filtered X = (∀ Y. finite Y ∧ Y ⊆ X → (∃ x ∈ X. ∀ y ∈ Y. x ≤ y))

```

```

definition downset-set :: 'a set ⇒ 'a set (⇓) where
  ⇓X = {y. ∃ x ∈ X. y ≤ x}

```

```

definition upset-set :: 'a set ⇒ 'a set (⇑) where
  ⇑X = {y. ∃ x ∈ X. x ≤ y}

```

```

definition downset :: 'a ⇒ 'a set (⇓) where
  ⇓ = ⇓ ∘ (λx. {x})

```

```

definition upset :: 'a ⇒ 'a set (⇑) where
  ⇑ = ⇑ ∘ (λx. {x})

```

```

definition downsets :: 'a set set where
  downsets = Fix ⇓

```

```

definition upsets :: 'a set set where
  upsets = Fix ⇑

```

abbreviation *downset-setp* $X \equiv X \in \text{downsets}$

abbreviation *upset-setp* $X \equiv X \in \text{upsets}$

definition *ideals* :: 'a set set **where**
 $\text{ideals} = \{X. X \neq \{\} \wedge \text{downset-setp } X \wedge \text{directed } X\}$

definition *filters* :: 'a set set **where**
 $\text{filters} = \{X. X \neq \{\} \wedge \text{upset-setp } X \wedge \text{filtered } X\}$

abbreviation *idealp* $X \equiv X \in \text{ideals}$

abbreviation *filterp* $X \equiv X \in \text{filters}$

end

abbreviation *Sup-pres* :: ('a::Sup \Rightarrow 'b::Sup) \Rightarrow bool **where**
 $\text{Sup-pres } f \equiv f \circ \text{Sup} = \text{Sup} \circ (\cdot) f$

abbreviation *Inf-pres* :: ('a::Inf \Rightarrow 'b::Inf) \Rightarrow bool **where**
 $\text{Inf-pres } f \equiv f \circ \text{Inf} = \text{Inf} \circ (\cdot) f$

abbreviation *sup-pres* :: ('a::sup \Rightarrow 'b::sup) \Rightarrow bool **where**
 $\text{sup-pres } f \equiv (\forall x y. f (x \sqcup y) = f x \sqcup f y)$

abbreviation *inf-pres* :: ('a::inf \Rightarrow 'b::inf) \Rightarrow bool **where**
 $\text{inf-pres } f \equiv (\forall x y. f (x \sqcap y) = f x \sqcap f y)$

abbreviation *bot-pres* :: ('a::bot \Rightarrow 'b::bot) \Rightarrow bool **where**
 $\text{bot-pres } f \equiv f \perp = \perp$

abbreviation *top-pres* :: ('a::top \Rightarrow 'b::top) \Rightarrow bool **where**
 $\text{top-pres } f \equiv f \top = \top$

abbreviation *Sup-dual* :: ('a::Sup \Rightarrow 'b::Inf) \Rightarrow bool **where**
 $\text{Sup-dual } f \equiv f \circ \text{Sup} = \text{Inf} \circ (\cdot) f$

abbreviation *Inf-dual* :: ('a::Inf \Rightarrow 'b::Sup) \Rightarrow bool **where**
 $\text{Inf-dual } f \equiv f \circ \text{Inf} = \text{Sup} \circ (\cdot) f$

abbreviation *sup-dual* :: ('a::sup \Rightarrow 'b::inf) \Rightarrow bool **where**
 $\text{sup-dual } f \equiv (\forall x y. f (x \sqcup y) = f x \sqcap f y)$

abbreviation *inf-dual* :: ('a::inf \Rightarrow 'b::sup) \Rightarrow bool **where**
 $\text{inf-dual } f \equiv (\forall x y. f (x \sqcap y) = f x \sqcup f y)$

abbreviation *bot-dual* :: ('a::bot \Rightarrow 'b::top) \Rightarrow bool **where**
 $\text{bot-dual } f \equiv f \perp = \top$

abbreviation *top-dual* :: ('a::top ⇒ 'b::bot) ⇒ bool **where**
top-dual f ≡ f ⊤ = ⊥

9.1 Duality via Locales

sublocale *ord* ⊆ *dual-ord*: *ord* (≥) (>
rewrites *dual-max-set*: *max-set* = *dual-ord.min-set*
and *dual-filtered*: *filtered* = *dual-ord.directed*
and *dual-upset-set*: *upset-set* = *dual-ord.downset-set*
and *dual-upset*: *upset* = *dual-ord.downset*
and *dual-upsets*: *upsets* = *dual-ord.downsets*
and *dual-filters*: *filters* = *dual-ord.ideals*
apply *unfold-locale*
unfolding *max-set-def* *ord.min-set-def* *fun-eq-iff* **apply** *blast*
unfolding *filtered-def* *ord.directed-def* **apply** *simp*
unfolding *upset-set-def* *ord.downset-set-def* **apply** *simp*
apply (*simp* *add*: *ord.downset-def* *ord.downset-set-def* *ord.upset-def* *ord.upset-set-def*)
unfolding *upsets-def* *ord.downsets-def* **apply** (*metis* *ord.downset-set-def* *up-set-set-def*)
unfolding *filters-def* *ord.ideals-def* *Fix-def* *ord.downsets-def* *upsets-def* *ord.downset-set-def*
upset-set-def *ord.directed-def* *filtered-def*
by *simp*

sublocale *preorder* ⊆ *dual-preorder*: *preorder* (≥) (>
apply *unfold-locale*
apply (*simp* *add*: *less-le-not-le*)
apply *simp*
using *order-trans* **by** *blast*

sublocale *order* ⊆ *dual-order*: *order* (≥) (>
by (*unfold-locale*, *simp*)

sublocale *lattice* ⊆ *dual-lattice*: *lattice* *sup* (≥) (>) *inf*
by (*unfold-locale*, *simp-all*)

sublocale *bounded-lattice* ⊆ *dual-bounded-lattice*: *bounded-lattice* *sup* (≥) (>) *inf*
⊤ ⊥
by (*unfold-locale*, *simp-all*)

sublocale *boolean-algebra* ⊆ *dual-boolean-algebra*: *boolean-algebra* λx y. x ⊔ −y
uminus *sup* (≥) (>) *inf* ⊤ ⊥
by (*unfold-locale*, *simp-all* *add*: *inf-sup-distrib1*)

sublocale *complete-lattice* ⊆ *dual-complete-lattice*: *complete-lattice* *Sup* *Inf* *sup* (≥)
(>) *inf* ⊤ ⊥
rewrites *dual-gfp*: *gfp* = *dual-complete-lattice.lfp*

proof –

show *class.complete-lattice* *Sup* *Inf* *sup* (≥) (>) *inf* ⊤ ⊥
by (*unfold-locale*, *simp-all* *add*: *Sup-upper* *Sup-least* *Inf-lower* *Inf-greatest*)

```

then interpret dual-complete-lattice: complete-lattice Sup Inf sup ( $\geq$ ) ( $>$ ) inf  $\top$ 
 $\perp$ .
show gfp = dual-complete-lattice.lfp
  unfolding gfp-def dual-complete-lattice.lfp-def fun-eq-iff by simp
qed

```

```

context ord
begin

```

```

lemma dual-min-set: min-set = dual-ord.max-set
  by (simp add: dual-ord.dual-max-set)

```

```

lemma dual-directed: directed = dual-ord.filtered
  by (simp add: dual-ord.dual-filtered)

```

```

lemma dual-downset: downset = dual-ord.upset
  by (simp add: dual-ord.dual-upset)

```

```

lemma dual-downset-set: downset-set = dual-ord.upset-set
  by (simp add: dual-ord.dual-upset-set)

```

```

lemma dual-downsets: downsets = dual-ord.upsets
  by (simp add: dual-ord.dual-upsets)

```

```

lemma dual-ideals: ideals = dual-ord.filters
  by (simp add: dual-ord.dual-filters)

```

```

end

```

```

context complete-lattice
begin

```

```

lemma dual-lfp: lfp = dual-complete-lattice.gfp
  by (simp add: dual-complete-lattice.dual-gfp)

```

```

end

```

9.2 Properties of Orderings, Again

```

context ord
begin

```

```

lemma directed-nonempty: directed  $X \implies X \neq \{\}$ 
  unfolding directed-def by fastforce

```

```

lemma directed-ub: directed  $X \implies (\forall x \in X. \forall y \in X. \exists z \in X. x \leq z \wedge y \leq z)$ 
  by (meson empty-subsetI directed-def finite.emptyI finite.insert insert-subset order-refl)

```

```

lemma downset-set-prop:  $\Downarrow = \text{Union} \circ (\cdot) \Downarrow$ 
  unfolding downset-set-def downset-def fun-eq-iff by fastforce

lemma downset-set-prop-var:  $\Downarrow X = (\bigcup x \in X. \Downarrow x)$ 
  by (simp add: downset-set-prop)

lemma downset-prop:  $\Downarrow x = \{y. y \leq x\}$ 
  unfolding downset-def downset-set-def fun-eq-iff comp-def by fastforce

end

context preorder
begin

lemma directed-prop:  $X \neq \{\}$   $\implies (\forall x \in X. \forall y \in X. \exists z \in X. x \leq z \wedge y \leq z)$ 
 $\implies$  directed  $X$ 
proof -
  assume h1:  $X \neq \{\}$ 
  and h2:  $\forall x \in X. \forall y \in X. \exists z \in X. x \leq z \wedge y \leq z$ 
  {fix  $Y$ 
  have finite  $Y \implies Y \subseteq X \implies (\exists x \in X. \forall y \in Y. y \leq x)$ 
  proof (induct rule: finite-induct)
    case empty
    then show ?case
      using h1 by blast
  next
    case (insert  $x$   $F$ )
    then show ?case
      by (metis h2 insert-iff insert-subset order-trans)
  qed}
  thus ?thesis
  by (simp add: directed-def)
qed

lemma directed-alt:  $\text{directed } X = (X \neq \{\} \wedge (\forall x \in X. \forall y \in X. \exists z \in X. x \leq z \wedge y \leq z))$ 
  by (metis directed-prop directed-nonempty directed-ub)

lemma downset-set-ext:  $\text{id} \leq \Downarrow$ 
  unfolding le-fun-def id-def downset-set-def by auto

lemma downset-set-iso: mono  $\Downarrow$ 
  unfolding mono-def downset-set-def by blast

lemma downset-set-idem [simp]:  $\Downarrow \circ \Downarrow = \Downarrow$ 
  unfolding fun-eq-iff downset-set-def comp-def using order-trans by auto

lemma downset-faithful:  $\Downarrow x \subseteq \Downarrow y \implies x \leq y$ 
  by (simp add: downset-prop subset-eq)

```

lemma *downset-iso-iff*: $(\downarrow x \subseteq \downarrow y) = (x \leq y)$
using *atMost-iff downset-prop order-trans* **by** *blast*

lemma *downset-directed-downset-var* [*simp*]: $\text{directed } (\downarrow X) = \text{directed } X$

proof

assume *h1*: *directed X*
{fix *Y*
assume *h2*: *finite Y* **and** *h3*: $Y \subseteq \downarrow X$
hence $\forall y. \exists x. y \in Y \longrightarrow x \in X \wedge y \leq x$
by (*force simp: downset-set-def*)
hence $\exists f. \forall y. y \in Y \longrightarrow f y \in X \wedge y \leq f y$
by (*rule choice*)
hence $\exists f. \text{finite } (f \text{ ` } Y) \wedge f \text{ ` } Y \subseteq X \wedge (\forall y \in Y. y \leq f y)$
by (*metis finite-imageI h2 image-subsetI*)
hence $\exists Z. \text{finite } Z \wedge Z \subseteq X \wedge (\forall y \in Y. \exists z \in Z. y \leq z)$
by *fastforce*
hence $\exists Z. \text{finite } Z \wedge Z \subseteq X \wedge (\forall y \in Y. \exists z \in Z. y \leq z) \wedge (\exists x \in X. \forall z \in Z. z \leq x)$
by (*metis directed-def h1*)
hence $\exists x \in X. \forall y \in Y. y \leq x$
by (*meson order-trans*)
thus *directed* $(\downarrow X)$
unfolding *directed-def downset-set-def* **by** *fastforce*
next
assume *directed* $(\downarrow X)$
thus *directed X*
unfolding *directed-def downset-set-def*
apply *clarsimp*
by (*smt (verit) Ball-Collect order-refl order-trans subsetCE*)
qed

lemma *downset-directed-downset* [*simp*]: $\text{directed } \circ \downarrow = \text{directed}$
unfolding *fun-eq-iff comp-def* **by** *simp*

lemma *directed-downset-ideals*: $\text{directed } (\downarrow X) = (\downarrow X \in \text{ideals})$

by (*metis (mono-tags, lifting) Fix-def comp-apply directed-alt downset-set-idem downsets-def ideals-def mem-Collect-eq*)

end

lemma *downset-iso*: $\text{mono } (\downarrow :: 'a :: \text{order} \Rightarrow 'a \text{ set})$
by (*simp add: downset-iso-iff mono-def*)

context *order*

begin

lemma *downset-inj*: *inj* \downarrow

by (*metis injI downset-iso-iff order.eq-iff*)

end

context *lattice*
begin

lemma *lat-ideals*: $X \in \text{ideals} = (X \neq \{\}) \wedge X \in \text{downsets} \wedge (\forall x \in X. \forall y \in X. x \sqcup y \in X)$

unfolding *ideals-def directed-alt downsets-def Fix-def downset-set-def*
using *local.sup.bounded-iff* **by** *blast*

end

context *bounded-lattice*
begin

lemma *bot-ideal*: $X \in \text{ideals} \implies \perp \in X$

unfolding *ideals-def downsets-def Fix-def downset-set-def* **by** *fastforce*

end

context *complete-lattice*
begin

lemma *Sup-downset-id* [*simp*]: $\text{Sup} \circ \downarrow = \text{id}$

using *Sup-atMost atMost-def downset-prop* **by** *fastforce*

lemma *downset-Sup-id*: $\text{id} \leq \downarrow \circ \text{Sup}$

by (*simp add: Sup-upper downset-prop le-funI subsetI*)

lemma *Inf-Sup-var*: $\bigsqcup (\bigcap x \in X. \downarrow x) = \bigsqcap X$

unfolding *downset-prop* **by** (*simp add: Collect-ball-eq Inf-eq-Sup*)

lemma *Inf-pres-downset-var*: $(\bigcap x \in X. \downarrow x) = \downarrow (\bigsqcap X)$

unfolding *downset-prop* **by** (*safe, simp-all add: le-Inf-iff*)

end

lemma *lfp-in-Fix*:

fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'a$

shows $\text{mono } f \implies \text{lfp } f \in \text{Fix } f$

using *Fix-def lfp-unfold* **by** *fastforce*

lemma *gfp-in-Fix*:

fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'a$

shows $\text{mono } f \implies \text{gfp } f \in \text{Fix } f$

using *Fix-def gfp-unfold* **by** *fastforce*

lemma *nonempty-Fix*:

```

fixes f :: 'a::complete-lattice => 'a
shows mono f ==> Fix f ≠ {}
using lfp-in-Fix by fastforce

```

9.3 Dual Properties of Orderings from Locales

These properties can be proved very smoothly overall. But only within the context of a class or locale!

```

context ord
begin

```

```

lemma filtered-nonempty: filtered X ==> X ≠ {}
  by (simp add: dual-filtered dual-ord.directed-nonempty)

```

```

lemma filtered-lb: filtered X ==> (∀ x ∈ X. ∀ y ∈ X. ∃ z ∈ X. z ≤ x ∧ z ≤ y)
  by (simp add: dual-filtered dual-ord.directed-ub)

```

```

lemma upset-set-prop: ↑ = Union ∘ (·) ↑
  by (simp add: dual-ord.downset-set-prop dual-upset dual-upset-set)

```

```

lemma upset-set-prop-var: ↑X = (∪ x ∈ X. ↑x)
  by (simp add: dual-ord.downset-set-prop-var dual-upset dual-upset-set)

```

```

lemma upset-prop: ↑x = {y. x ≤ y}
  by (simp add: dual-ord.downset-prop dual-upset)

```

```

end

```

```

context preorder
begin

```

```

lemma filtered-prop: X ≠ {} ==> (∀ x ∈ X. ∀ y ∈ X. ∃ z ∈ X. z ≤ x ∧ z ≤ y) ==>
  filtered X
  by (simp add: dual-filtered dual-preorder.directed-prop)

```

```

lemma filtered-alt: filtered X = (X ≠ {} ∧ (∀ x ∈ X. ∀ y ∈ X. ∃ z ∈ X. z ≤ x ∧
  z ≤ y))
  by (simp add: dual-filtered dual-preorder.directed-alt)

```

```

lemma upset-set-ext: id ≤ ↑
  by (simp add: dual-preorder.downset-set-ext dual-upset-set)

```

```

lemma upset-set-anti: mono ↑
  by (simp add: dual-preorder.downset-set-iso dual-upset-set)

```

```

lemma up-set-idem [simp]: ↑ ∘ ↑ = ↑
  by (simp add: dual-upset-set)

```

```

lemma upset-faithful: ↑x ⊆ ↑y ==> y ≤ x

```


by (*metis dual-preorder.downset-faithful dual-upset*)

lemma *upset-anti-iff*: $(\uparrow y \subseteq \uparrow x) = (x \leq y)$
by (*simp add: dual-preorder.downset-iso-iff dual-upset*)

lemma *upset-filtered-upset* [*simp*]: $\text{filtered} \circ \uparrow = \text{filtered}$
by (*simp add: dual-filtered dual-upset-set*)

lemma *filtered-upset-filters*: $\text{filtered} (\uparrow X) = (\uparrow X \in \text{filters})$
using *dual-filtered dual-preorder.directed-downset-ideals dual-upset-set ord.dual-filters*
by *fastforce*

end

context *order*
begin

lemma *upset-inj*: *inj* \uparrow
by (*simp add: dual-order.downset-inj dual-upset*)

end

context *lattice*
begin

lemma *lat-filters*: $X \in \text{filters} = (X \neq \{\}) \wedge X \in \text{upsets} \wedge (\forall x \in X. \forall y \in X. x \sqcap y \in X)$
by (*simp add: dual-filters dual-lattice.lat-ideals dual-ord.downsets-def dual-upset-set upsets-def*)

end

context *bounded-lattice*
begin

lemma *top-filter*: $X \in \text{filters} \implies \top \in X$
by (*simp add: dual-bounded-lattice.bot-ideal dual-filters*)

end

context *complete-lattice*
begin

lemma *Inf-upset-id* [*simp*]: $\text{Inf} \circ \uparrow = \text{id}$
by (*simp add: dual-upset*)

lemma *upset-Inf-id*: $\text{id} \leq \uparrow \circ \text{Inf}$
by (*simp add: dual-complete-lattice.downset-Sup-id dual-upset*)

lemma *Sup-Inf-var*: $\prod (\bigcap x \in X. \uparrow x) = \bigsqcup X$
 by (*simp add: dual-complete-lattice.Inf-Sup-var dual-upset*)

lemma *Sup-dual-upset-var*: $(\bigcap x \in X. \uparrow x) = \uparrow(\bigsqcup X)$
 by (*simp add: dual-complete-lattice.Inf-pres-downset-var dual-upset*)

end

9.4 Examples that Do Not Dualise

lemma *upset-anti*: *antimono* ($\uparrow :: 'a :: \text{order} \Rightarrow 'a \text{ set}$)
 by (*simp add: antimono-def upset-anti-iff*)

context *complete-lattice*
begin

lemma *fSup-unfold*: $(f :: \text{nat} \Rightarrow 'a) \ 0 \sqcup (\bigsqcup n. f (Suc\ n)) = (\bigsqcup n. f\ n)$
 apply (*intro order.antisym sup-least*)
 apply (*rule Sup-upper, force*)
 apply (*rule Sup-mono, force*)
 apply (*safe intro!: Sup-least*)
 by (*case-tac n, simp-all add: Sup-upper le-supI2*)

lemma *fInf-unfold*: $(f :: \text{nat} \Rightarrow 'a) \ 0 \sqcap (\prod n. f (Suc\ n)) = (\prod n. f\ n)$
 apply (*intro order.antisym inf-greatest*)
 apply (*rule Inf-greatest, safe*)
 apply (*case-tac n*)
 apply *simp-all*
 using *Inf-lower inf.coboundedI2* apply *force*
 apply (*simp add: Inf-lower*)
 by (*auto intro: Inf-mono*)

end

lemma *fun-isol*: $\text{mono } f \Longrightarrow \text{mono } ((\circ) f)$
 by (*simp add: le-fun-def mono-def*)

lemma *fun-isor*: $\text{mono } f \Longrightarrow \text{mono } (\lambda x. x \circ f)$
 by (*simp add: le-fun-def mono-def*)

lemma *Sup-sup-pres*:
 fixes $f :: 'a :: \text{complete-lattice} \Rightarrow 'b :: \text{complete-lattice}$
 shows $\text{Sup-pres } f \Longrightarrow \text{sup-pres } f$
 by (*metis (no-types, opaque-lifting) Sup-empty Sup-insert comp-apply image-insert sup-bot.right-neutral*)

lemma *Inf-inf-pres*:
 fixes $f :: 'a :: \text{complete-lattice} \Rightarrow 'b :: \text{complete-lattice}$

shows $\text{Inf-pres } f \implies \text{inf-pres } f$
by (*smt (verit) INF-insert comp-eq-elim dual-complete-lattice.Sup-empty dual-complete-lattice.Sup-insert inf-top.right-neutral*)

lemma *Sup-bot-pres*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{Sup-pres } f \implies \text{bot-pres } f$
by (*metis SUP-empty Sup-empty comp-eq-elim*)

lemma *Inf-top-pres*:
fixes $f :: 'a::\text{complete-lattice} \Rightarrow 'b::\text{complete-lattice}$
shows $\text{Inf-pres } f \implies \text{top-pres } f$
by (*metis INF-empty comp-eq-elim dual-complete-lattice.Sup-empty*)

context *complete-lattice*
begin

lemma *iso-Inf-subdistl*:
assumes $\text{mono } (f :: 'a \Rightarrow 'b::\text{complete-lattice})$
shows $f \circ \text{Inf} \leq \text{Inf} \circ (\cdot) f$
by (*simp add: assms complete-lattice-class.le-Inf-iff le-funI Inf-lower monoD*)

lemma *iso-Sup-supdistl*:
assumes $\text{mono } (f :: 'a \Rightarrow 'b::\text{complete-lattice})$
shows $\text{Sup} \circ (\cdot) f \leq f \circ \text{Sup}$
by (*simp add: assms complete-lattice-class.SUP-le-iff le-funI dual-complete-lattice.Inf-lower monoD*)

lemma *Inf-subdistl-iso*:
fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$
shows $f \circ \text{Inf} \leq \text{Inf} \circ (\cdot) f \implies \text{mono } f$
unfolding *mono-def le-fun-def comp-def* **by** (*metis complete-lattice-class.le-INF-iff Inf-atLeast atLeast-iff*)

lemma *Sup-supdistl-iso*:
fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$
shows $\text{Sup} \circ (\cdot) f \leq f \circ \text{Sup} \implies \text{mono } f$
unfolding *mono-def le-fun-def comp-def* **by** (*metis complete-lattice-class.SUP-le-iff Sup-atMost atMost-iff*)

lemma *supdistl-iso*:
fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$
shows $(\text{Sup} \circ (\cdot) f \leq f \circ \text{Sup}) = \text{mono } f$
using *Sup-supdistl-iso iso-Sup-supdistl* **by** *force*

lemma *subdistl-iso*:
fixes $f :: 'a \Rightarrow 'b::\text{complete-lattice}$
shows $(f \circ \text{Inf} \leq \text{Inf} \circ (\cdot) f) = \text{mono } f$
using *Inf-subdistl-iso iso-Inf-subdistl* **by** *force*

end

lemma *fSup-distr*: *Sup-pres* ($\lambda x. x \circ f$)
unfolding *fun-eq-iff comp-def*
by (*smt (verit) Inf.INF-cong SUP-apply Sup-apply*)

lemma *fSup-distr-var*: $\bigsqcup F \circ g = (\bigsqcup f \in F. f \circ g)$
unfolding *fun-eq-iff comp-def*
by (*smt (verit) Inf.INF-cong SUP-apply Sup-apply*)

lemma *fInf-distr*: *Inf-pres* ($\lambda x. x \circ f$)
unfolding *fun-eq-iff comp-def*
by (*smt (verit) INF-apply Inf.INF-cong Inf-apply*)

lemma *fInf-distr-var*: $\bigsqcap F \circ g = (\bigsqcap f \in F. f \circ g)$
unfolding *fun-eq-iff comp-def*
by (*smt (verit) INF-apply Inf.INF-cong Inf-apply*)

lemma *fSup-subdistl*:
assumes *mono* ($f :: 'a :: \text{complete-lattice} \Rightarrow 'b :: \text{complete-lattice}$)
shows $\text{Sup} \circ (\cdot) ((\circ) f) \leq (\circ) f \circ \text{Sup}$
using *assms* **by** (*simp add: SUP-least Sup-upper le-fun-def monoD image-comp*)

lemma *fSup-subdistl-var*:
fixes $f :: 'a :: \text{complete-lattice} \Rightarrow 'b :: \text{complete-lattice}$
shows $\text{mono } f \Longrightarrow (\bigsqcup g \in G. f \circ g) \leq f \circ \bigsqcup G$
by (*simp add: SUP-least Sup-upper le-fun-def monoD image-comp*)

lemma *fInf-subdistl*:
fixes $f :: 'a :: \text{complete-lattice} \Rightarrow 'b :: \text{complete-lattice}$
shows $\text{mono } f \Longrightarrow (\circ) f \circ \text{Inf} \leq \text{Inf} \circ (\cdot) ((\circ) f)$
by (*simp add: INF-greatest Inf-lower le-fun-def monoD image-comp*)

lemma *fInf-subdistl-var*:
fixes $f :: 'a :: \text{complete-lattice} \Rightarrow 'b :: \text{complete-lattice}$
shows $\text{mono } f \Longrightarrow f \circ \bigsqcap G \leq (\bigsqcap g \in G. f \circ g)$
by (*simp add: INF-greatest Inf-lower le-fun-def monoD image-comp*)

lemma *Inf-pres-downset*: *Inf-pres* ($\downarrow :: 'a :: \text{complete-lattice} \Rightarrow 'a \text{ set}$)
unfolding *downset-prop fun-eq-iff comp-def*
by (*safe, simp-all add: le-Inf-iff*)

lemma *Sup-dual-upset*: *Sup-dual* ($\uparrow :: 'a :: \text{complete-lattice} \Rightarrow 'a \text{ set}$)
unfolding *upset-prop fun-eq-iff comp-def*
by (*safe, simp-all add: Sup-le-iff*)

This approach could probably be combined with the explicit functor-based one. This may be good for proofs, but seems conceptually rather ugly.

end

10 Duality Based on a Data Type

```
theory Order-Lattice-Props-Wenzel
  imports Main
begin
```

```
unbundle lattice-syntax
```

10.1 Wenzel's Approach Revisited

This approach is similar to, but inferior to the explicit class-based one. The main caveat is that duality is not involutive with this approach, and this allows dualising less theorems.

I copy Wenzel's development [11] in this subsection and extend it with additional properties. I show only the most important properties.

```
datatype 'a dual = dual (un-dual: 'a) (∂)
```

```
notation un-dual (∂-)
```

```
lemma dual-inj: inj ∂
  using injI by fastforce
```

```
lemma dual-surj: surj ∂
  using dual.exhaust-sel by blast
```

```
lemma dual-bij: bij ∂
  by (simp add: bijI dual-inj dual-surj)
```

Dual is not idempotent, and I see no way of imposing this condition. Yet at least an inverse exists — namely un-dual..

```
lemma dual-inv1 [simp]: ∂- ∘ ∂ = id
  by fastforce
```

```
lemma dual-inv2 [simp]: ∂ ∘ ∂- = id
  by fastforce
```

```
lemma dual-inv-inj: inj ∂-
  by (simp add: dual.expand injI)
```

```
lemma dual-inv-surj: surj ∂-
  by (metis dual.sel surj-def)
```

```
lemma dual-inv-bij: bij ∂-
  by (simp add: bij-def dual-inv-inj dual-inv-surj)
```

lemma *dual-iff*: $(\partial x = y) \longleftrightarrow (x = \partial^- y)$
by *fastforce*

Isabelle data types come with a number of generic functions.

The functor `map-dual` lifts functions to dual types. Isabelle's generic definition is not straightforward to understand and use. Yet conceptually it can be explained as follows.

lemma *map-dual-def-var* [*simp*]: $(\text{map-dual}::('a \Rightarrow 'b) \Rightarrow 'a \text{ dual} \Rightarrow 'b \text{ dual}) f = \partial \circ f \circ \partial^-$
unfolding *fun-eq-iff comp-def* **by** $(\text{metis dual.map-sel dual-iff})$

lemma *map-dual-def-var2*: $\partial^- \circ \text{map-dual } f = f \circ \partial^-$
by $(\text{simp add: rewriteL-comp-comp})$

lemma *map-dual-func1*: $\text{map-dual } (f \circ g) = \text{map-dual } f \circ \text{map-dual } g$
unfolding *fun-eq-iff comp-def* **by** $(\text{metis dual.exhaust dual.map})$

lemma *map-dual-func2* : $\text{map-dual } id = id$
by *simp*

The functor `map-dual` has an inverse functor as well.

definition *map-dual-inv* :: $('a \text{ dual} \Rightarrow 'b \text{ dual}) \Rightarrow ('a \Rightarrow 'b)$ **where**
 $\text{map-dual-inv } f = \partial^- \circ f \circ \partial$

lemma *map-dual-inv-func1*: $\text{map-dual-inv } id = id$
by $(\text{simp add: map-dual-inv-def})$

lemma *map-dual-inv-func2*: $\text{map-dual-inv } (f \circ g) = \text{map-dual-inv } f \circ \text{map-dual-inv } g$
unfolding *fun-eq-iff comp-def map-dual-inv-def* **by** (metis dual-iff)

lemma *map-dual-inv1*: $\text{map-dual} \circ \text{map-dual-inv} = id$
unfolding *fun-eq-iff map-dual-def-var map-dual-inv-def comp-def id-def*
by (metis dual-iff)

lemma *map-dual-inv2*: $\text{map-dual-inv} \circ \text{map-dual} = id$
unfolding *fun-eq-iff map-dual-def-var map-dual-inv-def comp-def id-def*
by (metis dual-iff)

Hence dual is an isomorphism between categories.

lemma *subset-dual*: $(\partial ' X = Y) \longleftrightarrow (X = \partial^- ' Y)$
by $(\text{metis dual-inj image-comp image-inv-f-f inv-o-cancel dual-inv2})$

lemma *subset-dual1*: $(X \subseteq Y) \longleftrightarrow (\partial ' X \subseteq \partial ' Y)$
by $(\text{simp add: dual-inj inj-image-subset-iff})$

lemma *dual-ball*: $(\forall x \in X. P (\partial x)) \longleftrightarrow (\forall y \in \partial ' X. P y)$
by *simp*

lemma *dual-inv-ball*: $(\forall x \in X. P (\partial^- x)) \longleftrightarrow (\forall y \in \partial^- \text{ ` } X. P y)$
by *simp*

lemma *dual-all*: $(\forall x. P (\partial x)) \longleftrightarrow (\forall y. P y)$
by (*metis dual.collapse*)

lemma *dual-inv-all*: $(\forall x. P (\partial^- x)) \longleftrightarrow (\forall y. P y)$
by (*metis dual-inv-surj surj-def*)

lemma *dual-ex*: $(\exists x. P (\partial x)) \longleftrightarrow (\exists y. P y)$
by (*metis UNIV-I bex-imageD dual-surj*)

lemma *dual-inv-ex*: $(\exists x. P (\partial^- x)) \longleftrightarrow (\exists y. P y)$
by (*metis dual.sel*)

lemma *dual-Collect*: $\{\partial x \mid x. P (\partial x)\} = \{y. P y\}$
by (*metis dual.exhaust*)

lemma *dual-inv-Collect*: $\{\partial^- x \mid x. P (\partial^- x)\} = \{y. P y\}$
by (*metis dual.collapse dual.inject*)

lemma *fun-dual1*: $(f \circ \partial = g) \longleftrightarrow (f = g \circ \partial^-)$
by *auto*

lemma *fun-dual2*: $(\partial \circ f = g) \longleftrightarrow (f = \partial^- \circ g)$
by *auto*

lemma *fun-dual3*: $(f \circ (\cdot) \partial = g) \longleftrightarrow (f = g \circ (\cdot) \partial^-)$
unfolding *fun-eq-iff comp-def* **by** (*metis subset-dual*)

lemma *fun-dual4*: $(f = \partial^- \circ g \circ (\cdot) \partial) \longleftrightarrow (\partial \circ f \circ (\cdot) \partial^- = g)$
by (*metis fun-dual2 fun-dual3 o-assoc*)

The next facts show incrementally that the dual of a complete lattice is a complete lattice. This follows once again Wenzel.

instantiation *dual* :: (*ord*) *ord*
begin

definition *less-eq-dual-def*: $(\leq) = \text{rel-dual } (\geq)$

definition *less-dual-def*: $(<) = \text{rel-dual } (>)$

instance..

end

lemma *less-eq-dual-def-var*: $(x \leq y) = (\partial^- y \leq \partial^- x)$
apply (*rule antisym*)

```

apply (simp add: dual.rel-sel less-eq-dual-def)
by (simp add: dual.rel-sel less-eq-dual-def)

lemma less-dual-def-var:  $(x < y) = (\partial^- y < \partial^- x)$ 
by (simp add: dual.rel-sel less-dual-def)

instance dual :: (preorder) preorder
apply standard
apply (simp add: less-dual-def-var less-eq-dual-def-var less-le-not-le)
apply (simp add: less-eq-dual-def-var)
by (meson less-eq-dual-def-var order-trans)

instance dual :: (order) order
by (standard, simp add: dual.expand less-eq-dual-def-var)

lemma dual-anti:  $x \leq y \implies \partial y \leq \partial x$ 
by (simp add: dual-inj less-eq-dual-def the-inv-f-f)

lemma dual-anti-iff:  $(x \leq y) = (\partial y \leq \partial x)$ 
by (simp add: dual-inj less-eq-dual-def the-inv-f-f)

map-dual does not map isotone functions to antitone ones. It simply lifts
the type!

lemma mono f  $\implies$  mono (map-dual f)
unfolding map-dual-def-var mono-def by (metis comp-apply dual-anti less-eq-dual-def-var)

instantiation dual :: (lattice) lattice
begin

definition inf-dual-def:  $x \sqcap y = \partial (\partial^- x \sqcup \partial^- y)$ 

definition sup-dual-def:  $x \sqcup y = \partial (\partial^- x \sqcap \partial^- y)$ 

instance
by (standard, simp-all add: dual-inj inf-dual-def sup-dual-def less-eq-dual-def-var
the-inv-f-f)

end

instantiation dual :: (complete-lattice) complete-lattice
begin

definition Inf-dual-def:  $Inf = \partial \circ Sup \circ (\cdot) \partial^-$ 

definition Sup-dual-def:  $Sup = \partial \circ Inf \circ (\cdot) \partial^-$ 

definition bot-dual-def:  $\perp = \partial \top$ 

definition top-dual-def:  $\top = \partial \perp$ 

```


instance

by (*standard, simp-all add: Inf-dual-def top-dual-def Sup-dual-def bot-dual-def dual-inj le-INF-iff SUP-le-iff INF-lower SUP-upper less-eq-dual-def-var the-inv-f-f*)

end

Next, directed and filtered sets, upsets, downsets, filters and ideals in posets are defined.

context *ord*

begin

definition *directed* :: 'a set \Rightarrow bool **where**

directed $X = (\forall Y. \text{finite } Y \wedge Y \subseteq X \longrightarrow (\exists x \in X. \forall y \in Y. y \leq x))$

definition *filtered* :: 'a set \Rightarrow bool **where**

filtered $X = (\forall Y. \text{finite } Y \wedge Y \subseteq X \longrightarrow (\exists x \in X. \forall y \in Y. x \leq y))$

definition *downset-set* :: 'a set \Rightarrow 'a set (\Downarrow) **where**

$\Downarrow X = \{y. \exists x \in X. y \leq x\}$

definition *upset-set* :: 'a set \Rightarrow 'a set (\Uparrow) **where**

$\Uparrow X = \{y. \exists x \in X. x \leq y\}$

end

10.2 Examples that Do Not Dualise

Filtered and directed sets are dual.

Proofs could be simplified if dual was idempotent.

lemma *filtered-directed-dual*: *filtered* \circ (\cdot) $\partial =$ *directed*

proof –

{**fix** $X::$ 'a set

have (*filtered* \circ (\cdot) ∂) $X = (\forall Y. \text{finite } (\partial^- \text{ ' } Y) \wedge \partial^- \text{ ' } Y \subseteq X \longrightarrow (\exists x \in X. \forall y \in (\partial^- \text{ ' } Y). \partial x \leq \partial y))$

unfolding *filtered-def comp-def* **by** (*simp,metis dual-iff finite-subset-image subset-dual subset-dual1*)

also have ... = $(\forall Y. \text{finite } Y \wedge Y \subseteq X \longrightarrow (\exists x \in X. \forall y \in Y. y \leq x))$

by (*metis dual-anti-iff dual-inv-surj finite-subset-image top.extremum*)

finally have (*filtered* \circ (\cdot) ∂) $X =$ *directed* X

using *directed-def* **by** *auto*}

thus *?thesis*

unfolding *fun-eq-iff* **by** *simp*

qed

lemma *directed-filtered-dual*: *directed* \circ (\cdot) $\partial =$ *filtered*

proof –

{**fix** $X::$ 'a set

```

have (directed  $\circ$  ( $\cdot$ )  $\partial$ )  $X = (\forall Y. \text{finite } (\partial^- \text{ ' } Y) \wedge \partial^- \text{ ' } Y \subseteq X \longrightarrow (\exists x \in X. \forall y \in (\partial^- \text{ ' } Y). \partial y \leq \partial x))$ 
unfolding directed-def comp-def by (simp, metis dual-iff finite-subset-image subset-dual subset-dual1)
also have ... =  $(\forall Y. \text{finite } Y \wedge Y \subseteq X \longrightarrow (\exists x \in X. \forall y \in Y. x \leq y))$ 
unfolding dual-anti-iff[symmetric] by (metis dual-inv-surj finite-subset-image top-greatest)
finally have (directed  $\circ$  ( $\cdot$ )  $\partial$ )  $X = \text{filtered } X$ 
using filtered-def by auto
thus ?thesis
unfolding fun-eq-iff by simp
qed

```

This example illustrates the deficiency of the approach. In the class-based approach the second proof is trivial.

The next example shows that this is a systematic problem.

```

lemma downset-set-upset-set-dual: ( $\cdot$ )  $\partial \circ \Downarrow = \Uparrow \circ (\cdot) \partial$ 
proof –
  {fix  $X::'a \text{ set}$ 
have ( $(\cdot) \partial \circ \Downarrow$ )  $X = \{\partial y \mid y. \exists x \in X. y \leq x\}$ 
by (simp add: downset-set-def setcompr-eq-image)
also have ... =  $\{\partial y \mid y. \exists x \in X. \partial x \leq \partial y\}$ 
by (meson dual-anti-iff)
also have ... =  $\{y. \exists x \in \partial \text{ ' } X. x \leq y\}$ 
by (metis (mono-tags, opaque-lifting) dual.exhaust image-iff)
finally have ( $(\cdot) \partial \circ \Downarrow$ )  $X = (\Uparrow \circ (\cdot) \partial) X$ 
by (simp add: upset-set-def)}
thus ?thesis
unfolding fun-eq-iff by simp
qed

lemma upset-set-downset-set-dual: ( $\cdot$ )  $\partial \circ \Uparrow = \Downarrow \circ (\cdot) \partial$ 
unfolding downset-set-def upset-set-def fun-eq-iff comp-def
apply (safe, force simp: dual-anti)
by (metis (mono-tags, lifting) dual.exhaust dual-anti-iff mem-Collect-eq rev-image-eqI)

end

```

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