

No Faster-Than-Light Observers

Mike Stannett

April 14, 2026

Abstract

We provide a formal proof within First Order Relativity Theory that no observer can travel faster than the speed of light. Originally reported by Stannett and Némethi [1].

Contents

```
theory SpaceTime  
imports Main  
begin
```

```
record 'a Vector =  
  tdir :: 'a  
  xdir :: 'a  
  ydir :: 'a  
  zdir :: 'a
```

```
record 'a Point =  
  tval :: 'a  
  xval :: 'a  
  yval :: 'a  
  zval :: 'a
```

```
record 'a Line =  
  basepoint :: 'a Point  
  direction :: 'a Vector
```

```

record 'a Plane =
  pbasepoint :: 'a Point
  direction1  :: 'a Vector
  direction2  :: 'a Vector

```

```

record 'a Cone =
  vertex :: 'a Point
  slope  :: 'a

```

```

class Quantities = linordered-field

```

```

class Vectors = Quantities
begin

```

```

abbreviation vecZero :: 'a Vector (<0>) where
  vecZero ≡ (| tdir = (0::'a), xdir = 0, ydir = 0, zdir = 0 |)

```

```

fun vecPlus :: 'a Vector ⇒ 'a Vector ⇒ 'a Vector (infix <⊕> 100) where
  vecPlus u v = (| tdir = tdir u + tdir v, xdir = xdir u + xdir v,
    ydir = ydir u + ydir v, zdir = zdir u + zdir v |)

```

```

fun vecMinus :: 'a Vector ⇒ 'a Vector ⇒ 'a Vector (infix <⊖> 100) where
  vecMinus u v = (| tdir = tdir u - tdir v, xdir = xdir u - xdir v,
    ydir = ydir u - ydir v, zdir = zdir u - zdir v |)

```

```

fun vecNegate :: 'a Vector ⇒ 'a Vector (<~ ->) where
  vecNegate u = (| tdir = uminus (tdir u), xdir = uminus (xdir u),
    ydir = uminus (ydir u), zdir = uminus (zdir u) |)

```

```

fun innerProd :: 'a Vector ⇒ 'a Vector ⇒ 'a (infix <dot> 50) where
  innerProd u v = (tdir u * tdir v) + (xdir u * xdir v) +
    (ydir u * ydir v) + (zdir u * zdir v)

```

```

fun sqrlen :: 'a Vector ⇒ 'a where sqrlen u = (u dot u)

```

```

fun minkowskiProd :: 'a Vector ⇒ 'a Vector ⇒ 'a (infix <mdot> 50) where
  minkowskiProd u v = (tdir u * tdir v)
    - ((xdir u * xdir v) + (ydir u * ydir v) + (zdir u * zdir v))

```

```

fun mSqrLen :: 'a Vector ⇒ 'a where mSqrLen u = (u mdot u)

```

```

fun vecScale :: 'a ⇒ 'a Vector ⇒ 'a Vector (infix <*> 200) where

```

$vecScale\ k\ u = (\ tdir = k * tdir\ u, xdir = k * xdir\ u, ydir = k * ydir\ u, zdir = k * zdir\ u\)$

fun *orthogonal* :: 'a Vector \Rightarrow 'a Vector \Rightarrow bool (**infix** $\langle \perp \rangle$ 150) **where**
orthogonal *u v* = (*u dot v* = 0)

lemma *lemVecZeroMinus*:

shows $0 \ominus u = \sim u$

by *simp*

lemma *lemVecSelfMinus*:

shows $u \ominus u = 0$

by *simp*

lemma *lemVecPlusCommute*:

shows $u \oplus v = v \oplus u$

by (*simp add: add.commute*)

lemma *lemVecPlusAssoc*:

shows $u \oplus (v \oplus w) = (u \oplus v) \oplus w$

by (*simp add: add.assoc*)

lemma *lemVecPlusMinus*:

shows $u \oplus (\sim v) = u \ominus v$

by (*simp add: local.add-uminus-conv-diff*)

lemma *lemDotCommute*:

shows (*u dot v*) = (*v dot u*)

by (*simp add: mult.commute*)

lemma *lemMDotCommute*:

shows (*u mdot v*) = (*v mdot u*)

by (*simp add: mult.commute*)

lemma *lemScaleScale*:

shows $a**(b**u) = (a*b)**u$

by (*simp add: mult.assoc*)

lemma *lemScale1*:
 shows $1 ** u = u$
 by *simp*

lemma *lemScale0*:
 shows $0 ** u = 0$
 by *simp*

lemma *lemScaleNeg*:
 shows $(-k)**u = \sim (k**u)$
 by *simp*

lemma *lemScaleOrigin*:
 shows $k**0 = 0$
 by *auto*

lemma *lemScaleOverAdd*:
 shows $k**(u \oplus v) = k**u \oplus k**v$
 by (*simp add: semiring-normalization-rules(34)*)

lemma *lemAddOverScale*:
 shows $a**u \oplus b**u = (a+b)**u$
 by (*simp add: semiring-normalization-rules(1)*)

lemma *lemScaleInverse*:
 assumes $k \neq (0::'a)$
 and $v = k**u$
 shows $u = (\text{inverse } k)**v$
proof –
 have $(\text{inverse } k)**v = (\text{inverse } k * k)**u$
 by (*simp add: lemScaleScale assms(2) mult.assoc*)
 thus *?thesis* **by** (*metis (lifting) field-inverse assms(1) lemScale1*)
qed

```

lemma lemOrthoSym:
  assumes  $u \perp v$ 
  shows  $v \perp u$ 
  by (metis assms(1) lemDotCommute orthogonal.simps)

end

```

```

class Points = Quantities + Vectors
begin

```

```

  abbreviation origin :: 'a Point where
    origin  $\equiv$  ( $\mid$   $tval = 0, xval = 0, yval = 0, zval = 0$   $\mid$ )

```

```

  fun vectorJoining :: 'a Point  $\Rightarrow$  'a Point  $\Rightarrow$  'a Vector (from - to -) where
    vectorJoining  $p\ q$ 
      = ( $\mid$   $tdir = tval\ q - tval\ p, xdir = xval\ q - xval\ p,$ 
         $ydir = yval\ q - yval\ p, zdir = zval\ q - zval\ p$   $\mid$ )

```

```

  fun moveBy :: 'a Point  $\Rightarrow$  'a Vector  $\Rightarrow$  'a Point (infixl  $\langle \rightsquigarrow \rangle$  100) where
    moveBy  $p\ u$ 
      = ( $\mid$   $tval = tval\ p + tdir\ u, xval = xval\ p + xdir\ u,$ 
         $yval = yval\ p + ydir\ u, zval = zval\ p + zdir\ u$   $\mid$ )

```

```

  fun positionVector :: 'a Point  $\Rightarrow$  'a Vector where
    positionVector  $p =$  ( $\mid$   $tdir = tval\ p, xdir = xval\ p, ydir = yval\ p, zdir = zval\ p$   $\mid$ )

```

```

  fun before :: 'a Point  $\Rightarrow$  'a Point  $\Rightarrow$  bool (infixr  $\langle \lesssim \rangle$  100) where
    before  $p\ q = (tval\ p < tval\ q)$ 

```

```

  fun after :: 'a Point  $\Rightarrow$  'a Point  $\Rightarrow$  bool (infixr  $\langle \gtrsim \rangle$  100) where
    after  $p\ q = (tval\ p > tval\ q)$ 

```

```

  fun sametime :: 'a Point  $\Rightarrow$  'a Point  $\Rightarrow$  bool (infixr  $\langle \approx \rangle$  100) where
    sametime  $p\ q = (tval\ p = tval\ q)$ 

```

```

lemma lemFromToTo:

```

```

  shows (from  $p$  to  $q$ )  $\oplus$  (from  $q$  to  $r$ ) = (from  $p$  to  $r$ )

```

```

proof -

```

```

  have shared:  $\forall\ valp\ valq\ valr. (valq - valp + (valr - valq) = valr - valp)$ 

```

```

  by (metis add-uminus-conv-diff add-diff-cancel

```

```

    semiring-normalization-rules(24) semiring-normalization-rules(25))

```

```

  thus ?thesis by auto

```

```

qed

```

lemma *lemMoveByMove*:
shows $p \rightsquigarrow u \rightsquigarrow v = p \rightsquigarrow (u \oplus v)$
by (*simp add: add.assoc*)

lemma *lemScaleLinear*:
shows $p \rightsquigarrow a**u \rightsquigarrow b**v = p \rightsquigarrow (a**u \oplus b**v)$
by (*simp add: add.assoc*)

end

class *Lines* = *Quantities* + *Vectors* + *Points*
begin

fun *onAxisT* :: '*a* Point \Rightarrow bool **where**
onAxisT u = ((*xval* u = 0) \wedge (*yval* u = 0) \wedge (*zval* u = 0))

fun *space2* :: ('*a* Point) \Rightarrow ('*a* Point) \Rightarrow '*a* **where**
space2 u v
= (*xval* u - *xval* v) * (*xval* u - *xval* v)
+ (*yval* u - *yval* v) * (*yval* u - *yval* v)
+ (*zval* u - *zval* v) * (*zval* u - *zval* v)

fun *time2* :: ('*a* Point) \Rightarrow ('*a* Point) \Rightarrow '*a* **where**
time2 u v = (*tval* u - *tval* v) * (*tval* u - *tval* v)

fun *speed* :: ('*a* Point) \Rightarrow ('*a* Point) \Rightarrow '*a* **where**
speed u v = (*space2* u v / *time2* u v)

fun *mkLine* :: '*a* Point \Rightarrow '*a* Vector \Rightarrow '*a* Line **where**
mkLine b d = (\lfloor basepoint = b, direction = d \rfloor)

fun *lineJoining* :: '*a* Point \Rightarrow '*a* Point \Rightarrow '*a* Line (*line joining - to -*) **where**
lineJoining p q = (\lfloor basepoint = p, direction = from p to q \rfloor)

fun *parallel* :: '*a* Line \Rightarrow '*a* Line \Rightarrow bool (*parallel*) **where**
parallel lineA lineB = ((direction lineA = *vecZero*) \vee (direction lineB = *vecZero*)
 \vee (\exists k. (k \neq (0::'*a*) \wedge direction lineB = k**direction
lineA)))

fun *collinear* :: '*a* Point \Rightarrow '*a* Point \Rightarrow '*a* Point \Rightarrow bool **where**
collinear p q r = (\exists α β . ($\alpha + \beta = 1$) \wedge
positionVector p = α ** (positionVector q) \oplus β ** (positionVector r)))

fun *inLine* :: '*a* Point \Rightarrow '*a* Line \Rightarrow bool **where**
inLine p l = *collinear* p (basepoint l) (basepoint l \rightsquigarrow direction l)

```

fun meets :: 'a Line  $\Rightarrow$  'a Line  $\Rightarrow$  bool where
  meets line1 line2 = ( $\exists p.(inLine\ p\ line1 \wedge inLine\ p\ line2)$ )

```

```

lemma lemParallelReflexive:
  shows lineA  $\parallel$  lineA
proof -
  define dir where dir = direction lineA
  have (1  $\neq$  0)  $\wedge$  (dir = 1**dir) by simp
  thus ?thesis by (metis dir-def parallel.simps)
qed

```

```

lemma lemParallelSym:
  assumes lineA  $\parallel$  lineB
  shows lineB  $\parallel$  lineA
proof -
  have case1: direction lineA = vecZero  $\longrightarrow$  ?thesis by auto
  have case2: direction lineB = vecZero  $\longrightarrow$  ?thesis by auto
  {
    assume case3: direction lineA  $\neq$  vecZero  $\wedge$  direction lineB  $\neq$  vecZero
    have exists-kab:  $\exists kab.(kab \neq (0::'a) \wedge direction\ lineB = kab**direction\ lineA)$ 

      by (metis parallel.simps assms(1) case3)
      define kab where kab  $\equiv$  (SOME kab.(kab  $\neq$  (0::'a)  $\wedge$  direction lineB =
kab**direction lineA))
      have kab-props: kab  $\neq$  0  $\wedge$  direction lineB = kab**direction lineA
      using exists-kab kab-def
      by (rule Hilbert-Choice.exE-some)

      define kba where kba = inverse kab
      have kba-nonzero: kba  $\neq$  0 by (metis inverse-zero-imp-zero kab-props kba-def)
      have direction lineA = kba**direction lineB by (metis kba-def lemScaleInverse
kab-props)
      hence ?thesis by (metis kba-nonzero parallel.simps)
    }
  from this have (direction lineA  $\neq$  vecZero  $\wedge$  direction lineB  $\neq$  vecZero)  $\longrightarrow$ 
?thesis by blast

  thus ?thesis by (metis case1 case2)
qed

```

```

lemma lemParallelTrans:
  assumes lineA  $\parallel$  lineB
  and lineB  $\parallel$  lineC
  and direction lineB  $\neq$  vecZero
  shows lineA  $\parallel$  lineC
proof -

```

```

have case1: direction lineA = vecZero → ?thesis by auto
have case2: direction lineC = vecZero → ?thesis by auto

{
  assume case3: direction lineA ≠ vecZero ∧ direction lineC ≠ vecZero

  have exists-kab: ∃ kab.(kab ≠ (0::'a) ∧ direction lineB = kab**direction lineA)

    by (metis parallel.simps assms(1) case3 assms(3))
  then obtain kab where kab-props: kab ≠ 0 ∧ direction lineB = kab**direction
lineA by auto

  have exists-kbc: ∃ kbc.(kbc ≠ (0::'a) ∧ direction lineC = kbc**direction lineB)

    by (metis parallel.simps assms(2) case3 assms(3))
  then obtain kbc where kbc-props: kbc ≠ 0 ∧ direction lineC = kbc**direction
lineB by auto

  define kac where kac = kbc * kab
  have kac-nonzero: kac ≠ 0 by (metis kab-props kac-def kbc-props no-zero-divisors)
  have direction lineC = kac**direction lineA
    by (metis kab-props kbc-props kac-def lemScaleScale)
  hence ?thesis by (metis kac-nonzero parallel.simps)
}
from this have (direction lineA ≠ vecZero ∧ direction lineC ≠ vecZero) →
?thesis by blast

  thus ?thesis by (metis case1 case2)
qed

```

```

lemma (in -) lemLineIdentity:
  assumes lineA = ⟨ basepoint = basepoint lineB, direction = direction lineB ⟩
  shows lineA = lineB
proof -
have basepoint lineA = basepoint lineB ∧ direction lineA = direction lineB
  by (simp add: assms(1))
thus ?thesis by simp
qed

```

```

lemma lemDirectionJoining:
  shows vectorJoining p (p ~> v) = v
proof -
  have ∀ a b.(a + b - a = b)
  by (metis add-uminus-conv-diff diff-add-cancel semiring-normalization-rules(24))
  thus ?thesis by auto
qed

```

lemma *lemDirectionFromTo*:
shows *direction (line joining p to (p ~> dir)) = dir*
proof –
have *direction (line joining p to (p ~> dir)) = from p to (p ~> dir)* **by** *simp*
thus *?thesis* **by** (*metis lemDirectionJoining*)
qed

lemma *lemLineEndpoint*:
shows *q = p ~> (from p to q)*
proof –
have $\forall a b. (b = a + (b - a))$
by (*metis diff-add-cancel semiring-normalization-rules(24)*)
thus *?thesis* **by** *auto*
qed

lemma *lemNullLine*:
assumes *direction lineA = vecZero*
and *inLine x lineA*
shows *x = basepoint lineA*
proof –
define *bp* **where** *bp = basepoint lineA*
have *collinear x (basepoint lineA) (basepoint lineA ~> direction lineA)*
by (*metis inLine.simps assms(2)*)
hence *collinear x bp (bp ~> vecZero)* **by** (*metis bp-def assms(1)*)
hence *collinear x bp bp* **by** *simp*
hence $\exists a b. (a + b = 1) \wedge$
 $(\text{positionVector } x = a**(\text{positionVector } bp) \oplus b**(\text{positionVector } bp))$
by (*metis collinear.simps*)
hence *positionVector x = positionVector bp* **by** (*metis lemScale1 lemAddOverScale*)
thus *?thesis* **by** (*simp add: bp-def*)
qed

lemma *lemLineContainsBasepoint*:
shows *inLine p (line joining p to q)*
proof –
define *linePQ* **where** *linePQ = line joining p to q*
have *bp: basepoint linePQ = p* **by** (*simp add: linePQ-def*)
have *dir: direction linePQ = from p to q* **by** (*simp add: linePQ-def*)
have *endq: basepoint linePQ ~> direction linePQ = q* **by** (*metis bp dir lemLineEndpoint*)

have $(1 + 0 = 1) \wedge (\text{positionVector } p = 1**(\text{positionVector } p) \oplus 0**(\text{positionVector } q))$

by *auto*
hence *collinear p p q* **by** (*metis collinear.simps*)
hence *collinear p (basepoint linePQ) (basepoint linePQ ~ direction linePQ)*
by (*metis bp endq*)
thus *?thesis* **by** (*simp add: linePQ-def*)
qed

lemma *lemLineContainsEndpoint*:
shows *inLine q (line joining p to q)*
proof –
define *linePQ* **where** *linePQ = line joining p to q*
have *bp: basepoint linePQ = p* **by** (*simp add: linePQ-def*)
have *dir: direction linePQ = from p to q* **by** (*simp add: linePQ-def*)
have *endq: basepoint linePQ ~ direction linePQ = q* **by** (*metis bp dir lemLineEndpoint*)

have ($0 + 1 = 1$) \wedge (*positionVector q = 0** (positionVector p) \oplus 1** (positionVector q)*)
by *auto*
hence *collinear q p q* **by** (*metis collinear.simps*)
hence *collinear q (basepoint linePQ) (basepoint linePQ ~ direction linePQ)*
by (*metis bp endq*)
thus *?thesis* **by** (*simp add: linePQ-def*)
qed

lemma *lemDirectionReverse*:
shows *from q to p = vecNegate (from p to q)*
by *simp*

lemma *lemParallelJoin*:
assumes *line joining p to q || line joining q to r*
shows *line joining p to q || line joining p to r*
proof –
define *linePQ* **where** *linePQ = line joining p to q*
define *lineQR* **where** *lineQR = line joining q to r*
define *linePR* **where** *linePR = line joining p to r*

have *case1: (direction linePQ = vecZero) \longrightarrow ?thesis* **by** (*simp add: linePQ-def*)
have *case2: (direction linePR = vecZero) \longrightarrow ?thesis* **by** (*simp add: linePR-def*)

{
assume *case3: direction linePQ \neq vecZero \wedge direction linePR \neq vecZero*
{
assume *case3a: direction lineQR = vecZero*
have *inLine r lineQR* **by** (*metis lemLineContainsEndpoint lineQR-def*)
}
}

```

    hence  $r = \text{basepoint } \text{lineQR}$  by (metis lemNullLine case3a)
    hence  $r = q$  by (simp add: lineQR-def)
    hence  $\text{linePQ} = \text{linePR}$  by (simp add: linePQ-def linePR-def)
    hence ?thesis by (metis lemParallelReflexive linePQ-def linePR-def)
  }
  from this have rtp3a:  $\text{direction } \text{lineQR} = \text{vecZero} \longrightarrow ?thesis$  by blast

  {
    assume case3b:  $\text{direction } \text{lineQR} \neq \text{vecZero}$ 

    define dirPQ where  $\text{dirPQ} = \text{from } p \text{ to } q$ 
    have dir-pq:  $\text{direction } \text{linePQ} = \text{dirPQ}$  by (simp add: linePQ-def dirPQ-def)

    define dirQR where  $\text{dirQR} = \text{from } q \text{ to } r$ 
    have dir-qr:  $\text{direction } \text{lineQR} = \text{dirQR}$  by (simp add: lineQR-def dirQR-def)

    have exists-k:  $\exists k. (k \neq 0 \wedge \text{direction } \text{lineQR} = k ** \text{direction } \text{linePQ})$ 
      by (metis linePQ-def lineQR-def assms(1) parallel.simps case3b case3)
    then obtain k where k-props:  $k \neq 0 \wedge \text{dirQR} = k ** \text{dirPQ}$  by (metis dir-pq
dir-qr)

    define scalar where  $\text{scalar} = 1 + k$ 

    have  $q = p \rightsquigarrow \text{dirPQ} \wedge r = q \rightsquigarrow \text{dirQR}$  by (metis lemLineEndpoint
dirPQ-def dirQR-def)
    hence  $r = p \rightsquigarrow \text{dirPQ} \rightsquigarrow (k ** \text{dirPQ})$  by (metis k-props)
    hence scalarPR:  $r = p \rightsquigarrow \text{scalar} ** \text{dirPQ}$ 
      by (metis lemScaleLinear lemScale1 lemAddOverScale scalar-def)

    {
      assume scalar0:  $\text{scalar} = 0$ 
      have  $r = p$  by (simp add: lemScale0 scalarPR scalar0)
      hence  $\text{direction } \text{linePR} = \text{vecZero}$  by (simp add: linePR-def)
      hence False by (metis case3)
    }
    from this have scalar-nonzero:  $\text{scalar} \neq 0$  by blast

    have  $\text{linePR} = \text{line joining } p \text{ to } (p \rightsquigarrow \text{scalar} ** \text{dirPQ})$ 
      by (simp add: linePR-def scalarPR)
    hence  $\text{direction } \text{linePR} = \text{scalar} ** \text{dirPQ}$  by (metis lemDirectionFromTo)

    hence scalar-props:  $\text{scalar} \neq 0 \wedge \text{direction } \text{linePR} = \text{scalar} ** \text{direction } \text{linePQ}$ 
      by (metis scalar-nonzero dir-pq)
    hence ?thesis by (metis parallel.simps linePR-def linePQ-def)
  }
  from this have  $\text{direction } \text{lineQR} \neq \text{vecZero} \longrightarrow ?thesis$  by blast

  hence ?thesis by (metis rtp3a)

```

}
from this have ($\text{direction linePQ} \neq \text{vecZero} \wedge \text{direction linePR} \neq \text{vecZero}$) \longrightarrow
?thesis by blast

thus ?thesis by (*metis case1 case2*)
qed

lemma *lemDirectionCollinear*:

shows $\text{collinear } u \ v \ (v \rightsquigarrow d) \longleftrightarrow (\exists \beta. (\text{from } u \text{ to } v = (-\beta)**d))$

proof –

have *basic1*: $\forall u \ v. (\text{positionVector } (u \rightsquigarrow v)) = (\text{positionVector } u) \oplus v$ **by** *simp*

have *basic2*: $\forall u \ v \ w. (u = v \oplus w \longrightarrow v \ominus u = \text{vecNegate } w)$

apply *auto*

by (*metis add-uminus-conv-diff diff-add-cancel minus-add*
semiring-normalization-rules(24)) +

have *basic3*: $\forall u \ v. (\text{from } u \text{ to } v = \text{positionVector } v \ominus \text{positionVector } u)$ **by** *simp*

have *basic4*: $\forall u \ v \ w. (v \ominus u = \text{vecNegate } w \longrightarrow u = v \oplus w)$

apply *auto*

by (*metis add-uminus-conv-diff diff-add-cancel lemScale1 mult.left-neutral*
semiring-normalization-rules(24) vecScale.simps)

{

assume *assm*: $\text{collinear } u \ v \ (v \rightsquigarrow d)$

have $\exists \alpha \ \beta. (\alpha + \beta = 1) \wedge$

$\text{positionVector } u = \alpha**(\text{positionVector } v) \oplus \beta**(\text{positionVector } (v \rightsquigarrow d))$

by (*metis assm collinear.simps*)

then obtain $\alpha \ \beta$ **where** *props*: $(\alpha + \beta = 1) \wedge$

$\text{positionVector } u = \alpha**(\text{positionVector } v) \oplus \beta**(\text{positionVector } (v \rightsquigarrow$

$d))$ **by** *auto*

hence $\text{positionVector } u = 1**(\text{positionVector } v) \oplus \beta**d$

by (*metis basic1 lemScaleOverAdd lemVecPlusAssoc lemAddOverScale props*)

hence $\text{positionVector } u = \text{positionVector } v \oplus \beta**d$ **by** (*metis lemScale1*)

hence $\text{positionVector } v \ominus \text{positionVector } u = (-\beta)**d$ **by** (*metis basic2*

lemScaleNeg)

hence $\exists \beta. (\text{from } u \text{ to } v = (-\beta)**d)$ **by** (*metis basic3*)

}

from this have *fwd*: $\text{collinear } u \ v \ (v \rightsquigarrow d) \longrightarrow (\exists \beta. (\text{from } u \text{ to } v = (-\beta)**d))$

by *blast*

{

assume $\exists \beta. (\text{from } u \text{ to } v = (-\beta)**d)$

then obtain β **where** *assm*: $\text{from } u \text{ to } v = (-\beta)**d$ **by** *auto*

define α **where** $\alpha = 1 - \beta$

have $\alpha\beta\text{-sum}$: $\alpha + \beta = 1$ **by** (*simp add: $\alpha\text{-def}$*)

have $\text{from } u \text{ to } v = \text{vecNegate } (\beta**d)$ **by** (*metis assm lemScaleNeg*)

hence $\text{positionVector } v \ominus \text{positionVector } u = \text{vecNegate } (\beta**d)$ **by** *auto*

hence $\text{positionVector } u = \text{positionVector } v \oplus \beta**d$ **by** (*metis basic4*)

hence $\text{positionVector } u = 1**(\text{positionVector } v) \oplus \beta**d$

by (*metis lemScale1*)
hence $(\alpha + \beta = 1) \wedge$
 $\text{positionVector } u = \alpha**(\text{positionVector } v) \oplus \beta**(\text{positionVector } (v \rightsquigarrow d))$
by (*metis $\alpha\beta$ -sum basic1 lemScaleOverAdd lemVecPlusAssoc lemAddOverScale*)
hence *collinear u v (v \rightsquigarrow d)* **by** *auto*
}
from this have $(\exists \beta.(\text{from } u \text{ to } v = (-\beta)**d)) \longrightarrow \text{collinear } u \text{ v } (v \rightsquigarrow d)$ **by**
blast

thus *?thesis* **by** (*metis fwd*)
qed

lemma *lemParallelNotMeet:*

assumes *lineA* \parallel *lineB*
and *direction lineA* \neq *vecZero*
and *direction lineB* \neq *vecZero*
and *inLine x lineA*
and $\neg(\text{inLine } x \text{ lineB})$
shows $\neg(\text{meets } \text{lineA } \text{lineB})$
proof –

have *basic:* $\forall p \ q \ v \ a.(\text{from } p \text{ to } q = a**v \longrightarrow \text{from } q \text{ to } p = (-a)**v)$
apply (*simp add: lemScaleNeg*) **by** (*metis minus-diff-eq*)

define *bpA* **where** *bpA* = *basepoint lineA*
define *dirA* **where** *dirA* = *direction lineA*
define *bpB* **where** *bpB* = *basepoint lineB*
define *dirB* **where** *dirB* = *direction lineB*

have *lineB* \parallel *lineA* **by** (*metis lemParallelSym assms(1)*)
hence *exists-kab:* $\exists \text{ kab} . (\text{ kab } \neq (0::'a) \wedge \text{ direction } \text{lineA} = \text{ kab}**\text{direction } \text{lineB})$

by (*metis parallel.simps assms(2) assms(3)*)
then obtain *kab* **where** *kab-props:* $\text{ kab } \neq 0 \wedge \text{ dirA} = \text{ kab}**\text{dirB}$ **by** (*metis dirA-def dirB-def*)

have *collinear x bpA (bpA \rightsquigarrow dirA)* **by** (*metis assms(4) inLine.simps bpA-def dirA-def*)

then obtain β **where** *from x to bpA* = $(-\beta)**\text{dirA}$ **by** (*metis lemDirectionCollinear*)

hence *x-to-bpA:* *from x to bpA* = $((-\beta)*\text{ kab})**\text{dirB}$ **by** (*metis lemScaleScale kab-props*)

{

assume *converse: meets lineA lineB*
have $\exists p.(inLine\ p\ lineA \wedge inLine\ p\ lineB)$ **by** (*metis converse meets.simps*)
then obtain p **where** *p-in-AB: inLine p lineA \wedge inLine p lineB* **by** *auto*

have *collinear p bpA (bpA \rightsquigarrow dirA)* **by** (*metis p-in-AB inLine.simps bpA-def dirA-def*)
then obtain βA **where** *from p to bpA = $(-\beta A)**dirA$* **by** (*metis lemDirectionCollinear*)
hence *from bpA to p = $(\beta A)**dirA$* **by** (*metis basic minus-minus*)
hence *bpA-to-p: from bpA to p = $(\beta A*kab)**dirB$* **by** (*metis lemScaleScale kab-props*)

have *collinear p bpB (bpB \rightsquigarrow dirB)* **by** (*metis p-in-AB inLine.simps bpB-def dirB-def*)
then obtain βB **where** *p-to-bpB: from p to bpB = $(-\beta B)**dirB$* **by** (*metis lemDirectionCollinear*)

define γ **where** $\gamma = -((-\beta)*kab + (\beta A*kab) + (-\beta B))$
have *x-to-bpB: (from x to bpA) \oplus (from bpA to p) \oplus (from p to bpB) = (from x to bpB)*
by (*metis lemFromToTo*)
hence *from x to bpB = $((-\beta)*kab)**dirB \oplus (\beta A*kab)**dirB \oplus (-\beta B)**dirB$*
by (*metis x-to-bpA bpA-to-p p-to-bpB*)
hence *from x to bpB = $(-\gamma)**dirB$*
by (*metis lemAddOverScale add.assoc γ -def minus-minus*)
hence *collinear x bpB (bpB \rightsquigarrow dirB)* **by** (*metis lemDirectionCollinear*)
hence *inLine x lineB* **by** (*metis inLine.simps bpB-def dirB-def*)
}
from this have *meets lineA lineB \longrightarrow inLine x lineB* **by** *blast*
thus *?thesis* **by** (*metis assms(5)*)
qed

lemma *lemAxisIsLine:*

assumes *onAxisT x*
and *onAxisT y*
and *onAxisT z*
and $x \neq y$
and $y \neq z$
and $z \neq x$
shows *collinear x y z*

proof –

define *ratio* **where** $ratio = -(tval\ y - tval\ x) / (tval\ z - tval\ y)$

have *x-onAxis: xval x = 0 \wedge yval x = 0 \wedge zval x = 0* **by** (*metis assms(1) onAxisT.simps*)

have *y-onAxis: xval y = 0 \wedge yval y = 0 \wedge zval y = 0* **by** (*metis assms(2)*)

```

onAxisT.simps)
  have z-onAxis: xval z = 0 ∧ yval z = 0 ∧ zval z = 0 by (metis assms(3)
onAxisT.simps)

  have tval z - tval y = 0 → z = y by (simp add: z-onAxis y-onAxis)
  hence tval z ≠ tval y by (metis assms(5) eq-iff-diff-eq-0)
  hence tvalyz-nonzero: tval z - tval y ≠ 0 by (metis eq-iff-diff-eq-0)

  have x-to-y: from x to y = (| tdir = tval y - tval x, xdir = 0, ydir = 0, zdir
= 0 |)
  by (simp add: x-onAxis y-onAxis)
  have y-to-z: from y to z = (| tdir = tval z - tval y, xdir = 0, ydir = 0, zdir
= 0 |)
  by (simp add: y-onAxis z-onAxis)

  have from x to y = (-ratio)**(from y to z)
  apply (simp add: x-to-y y-to-z ratio-def)
  by (metis diff-self eq-divide-imp minus-diff-eq mult-eq-0-iff
tvalyz-nonzero x-onAxis y-onAxis z-onAxis)
  hence collinear x y (y ~ (from y to z)) by (metis lemDirectionCollinear)
  thus ?thesis by (metis lemLineEndpoint)
qed

lemma lemSpace2Sym:
  shows space2 x y = space2 y x
proof -
  define xsep where xsep = xval x - xval y
  define ysep where ysep = yval x - yval y
  define zsep where zsep = zval x - zval y

  have spacexy: space2 x y = (xsep*xsep) + (ysep*ysep) + (zsep*zsep)
  by (simp add: xsep-def ysep-def zsep-def)
  have spaceyx: space2 y x = (-xsep)*(-xsep) + (-ysep)*(-ysep) + (-zsep)*(-zsep)
  by (simp add: xsep-def ysep-def zsep-def)
  thus ?thesis by (metis spacexy diff-0-right minus-diff-eq minus-mult-left mi-
nus-mult-right)
qed

lemma lemTime2Sym:
  shows time2 x y = time2 y x
proof -
  define tsep where tsep = tval x - tval y

  have timexy: time2 x y = tsep*tsep
  by (simp add: tsep-def)
  have timeyx: time2 y x = (-tsep)*(-tsep)
  by (simp add: tsep-def)
  thus ?thesis by (metis timexy diff-0-right minus-diff-eq minus-mult-left mi-
nus-mult-right)

```

```

qed

end

class Planes = Quantities + Lines
begin
  fun mkPlane :: 'a Point ⇒ 'a Vector ⇒ 'a Vector ⇒ 'a Plane where
    mkPlane b d1 d2 = (| pbasepoint = b, direction1 = d1, direction2 = d2 |)

  fun coplanar :: 'a Point ⇒ 'a Point ⇒ 'a Point ⇒ 'a Point ⇒ bool where
    coplanar e x y z
      = (∃ α β γ. (α + β + γ = 1) ∧
        positionVector e
          = (α**(positionVector x) ⊕ β**(positionVector y) ⊕ γ**(positionVector
z) )))

  fun inPlane :: 'a Point ⇒ 'a Plane ⇒ bool where
    inPlane e pl = coplanar e (pbasepoint pl) (pbasepoint pl ~ direction1 pl)
      (pbasepoint pl ~ direction2 pl)

  fun samePlane :: 'a Plane ⇒ 'a Plane ⇒ bool where
    samePlane pl pl' = (inPlane (pbasepoint pl) pl' ∧
      inPlane (pbasepoint pl ~ direction1 pl) pl' ∧
      inPlane (pbasepoint pl ~ direction2 pl) pl')

lemma lemPlaneContainsBasePoint:
shows inPlane (pbasepoint pl) pl
proof –
  define α where α = (1::'a)
  define β where β = (0::'a)
  define γ where γ = (0::'a)
  have rtp1: α + β + γ = 1 by (simp add: α-def β-def γ-def)

  define e where e = pbasepoint pl
  define x where x = pbasepoint pl
  define y where y = pbasepoint pl ~ direction1 pl
  define z where z = pbasepoint pl ~ direction2 pl
  have rtp2: positionVector e = α**(positionVector x)
    ⊕ β**(positionVector y) ⊕ γ**(positionVector z)
    by (simp add: e-def x-def α-def β-def γ-def)

  have sameplane: coplanar e x y z by (metis coplanar.simps rtp1 rtp2)
  hence coplanar e (pbasepoint pl) (pbasepoint pl ~ direction1 pl)
    (pbasepoint pl ~ direction2 pl)
    by (simp add: x-def y-def z-def)
  hence inPlane e pl by simp
  thus ?thesis by (simp add: e-def)
qed

```

end

class *Cones* = *Quantities* + *Lines* + *Planes* +
fixes

tangentPlane :: 'a *Point* \Rightarrow 'a *Cone* \Rightarrow 'a *Plane*

assumes

AxTangentBase: *pbasepoint* (*tangentPlane* *e cone*) = *e*

and

AxTangentVertex: *inPlane* (*vertex cone*) (*tangentPlane* *e cone*)

and

AxConeTangent: (*onCone* *e cone*) \longrightarrow
 $((\text{inPlane } pt \text{ (tangentPlane } e \text{ cone)} \wedge \text{onCone } pt \text{ cone})$
 $\longleftrightarrow \text{collinear (vertex cone) } e \text{ } pt)$

and

AxParallelCones: (*onCone* *e econe* \wedge *e* \neq *vertex econe* \wedge *onCone* *f fccone* \wedge *f* \neq
vertex fccone

$\wedge \text{inPlane } f \text{ (tangentPlane } e \text{ econe)})$
 $\longrightarrow (\text{samePlane (tangentPlane } e \text{ econe) (tangentPlane } f \text{ fccone)}$
 $\wedge ((\text{lineJoining (vertex econe) } e) \parallel (\text{lineJoining (vertex fccone}$

f)))

and

AxParallelConesE: *outsideCone* *f cone*

$\longrightarrow (\exists e. (\text{onCone } e \text{ cone} \wedge e \neq \text{vertex cone} \wedge \text{inPlane } f \text{ (tangentPlane } e \text{ cone)}))$

and

AxSlopedLineInVerticalPlane: $\llbracket \text{onAxisT } e; \text{onAxisT } f; e \neq f; \neg(\text{onAxisT } g) \rrbracket$

$\implies (\forall s. (\exists p. (\text{collinear } e \text{ } g \text{ } p \wedge (\text{space2 } p \text{ } f = (s*s)*\text{time2 } p \text{ } f))))$

begin

fun *onCone* :: 'a *Point* \Rightarrow 'a *Cone* \Rightarrow *bool* **where**

onCone *p cone*

= (*space2* (*vertex cone*) *p* = (*slope cone* * *slope cone*) * *time2* (*vertex cone*)

p)

fun *insideCone* :: 'a *Point* \Rightarrow 'a *Cone* \Rightarrow *bool* **where**

insideCone *p cone*

= (*space2* (*vertex cone*) *p* < (*slope cone* * *slope cone*) * *time2* (*vertex cone*)

p)

```

fun outsideCone :: 'a Point  $\Rightarrow$  'a Cone  $\Rightarrow$  bool where
  outsideCone p cone
    = (space2 (vertex cone) p > (slope cone * slope cone) * time2 (vertex cone)
p)

```

```

fun mkCone :: 'a Point  $\Rightarrow$  'a  $\Rightarrow$  'a Cone where
  mkCone v s = (| vertex = v, slope = s |)

```

```

lemma lemVertexOnCone:
  shows onCone (vertex cone) cone
by simp

```

```

lemma lemOutsideNotOnCone:
  assumes outsideCone f cone
  shows  $\neg$  (onCone f cone)
by (metis assms less-irrefl onCone.simps outsideCone.simps)

```

end

```

class SpaceTime = Quantities + Vectors + Points + Lines + Planes + Cones

```

end

```

theory SomeFunc
  imports Main
begin

```

```

fun someFunc :: ('a  $\Rightarrow$  'b  $\Rightarrow$  bool)  $\Rightarrow$  'a  $\Rightarrow$  'b where
  someFunc P x = (SOME y. (P x y))

```

```

lemma lemSomeFunc:
  assumes  $\exists y . P x y$ 
  and f = someFunc P
  shows P x (f x)
proof -
  have f x = (SOME y. (P x y))
  using assms(2) by simp
  thus ?thesis using assms(1)
  by (simp add: someI-ex)
qed

```

end

```

theory Axioms

```

```

imports SpaceTime SomeFunc
begin

record Body =
  Ph :: bool
  IOb :: bool

class WorldView = SpaceTime +
fixes

  W :: Body  $\Rightarrow$  Body  $\Rightarrow$  'a Point  $\Rightarrow$  bool ( $\leftarrow$  sees - at  $\rightarrow$ )
and

  wvt :: Body  $\Rightarrow$  Body  $\Rightarrow$  'a Point  $\Rightarrow$  'a Point
assumes
  AxWVT:  $\llbracket$  IOb m; IOb k  $\rrbracket \Longrightarrow (W\ k\ b\ x \longleftrightarrow W\ m\ b\ (wvt\ m\ k\ x))$ 
and
  AxWVTSym:  $\llbracket$  IOb m; IOb k  $\rrbracket \Longrightarrow (y = wvt\ k\ m\ x \longleftrightarrow x = wvt\ m\ k\ y)$ 
begin
end

class AxiomPreds = WorldView
begin
  fun sqrtTest :: 'a  $\Rightarrow$  'a  $\Rightarrow$  bool where
    sqrtTest x r =  $((r \geq 0) \wedge (r*r = x))$ 

  fun cTest :: Body  $\Rightarrow$  'a  $\Rightarrow$  bool where
    cTest m v =  $( (v > 0) \wedge ( \forall x\ y . ($ 
       $(\exists p. (Ph\ p \wedge W\ m\ p\ x \wedge W\ m\ p\ y)) \longleftrightarrow (space2\ x\ y = (v * v)*(time2$ 
x y))
       $)))$ 
end

class AxEuclidean = AxiomPreds + Quantities +
assumes
  AxEuclidean:  $(x \geq Groups.zero-class.zero) \Longrightarrow (\exists r. sqrtTest\ x\ r)$ 
begin

  abbreviation sqrt :: 'a  $\Rightarrow$  'a where
    sqrt  $\equiv$  someFunc sqrtTest

  lemma lemSqrt:

```

```

assumes  $x \geq 0$ 
and  $r = \text{sqrt } x$ 
shows  $r \geq 0 \wedge r*r = x$ 
proof -
have  $\text{rootExists}: (\exists r. \text{sqrtTest } x \ r)$  by ( $\text{metis AxEuclidean assms}(1)$ )
hence  $\text{sqrtTest } x \ (\text{sqrt } x)$  by ( $\text{metis lemSomeFunc}$ )
thus  $?thesis$  using  $\text{assms}(2)$  by  $\text{simp}$ 
qed

end

class AxLight = WorldView +
assumes
   $\text{AxLight}: \exists m \ v. ( \text{IOb } m \wedge (v > (0::'a)) \wedge ( \forall x \ y. ($ 
     $(\exists p. (\text{Ph } p \wedge \text{W } m \ p \ x \wedge \text{W } m \ p \ y)) \longleftrightarrow (\text{space2 } x \ y = (v * v) * \text{time2 } x$ 
   $y))$ 
begin
end

class AxPh = WorldView + AxiomPreds +
assumes
   $\text{AxPh}: \text{IOb } m \implies (\exists v. \text{cTest } m \ v)$ 
begin

  abbreviation  $c :: \text{Body} \Rightarrow 'a$  where
     $c \equiv \text{someFunc } \text{cTest}$ 

  fun  $\text{lightcone} :: \text{Body} \Rightarrow 'a \ \text{Point} \Rightarrow 'a \ \text{Cone}$  where
     $\text{lightcone } m \ v = \text{mkCone } v \ (c \ m)$ 

lemma lemCProps:
assumes  $\text{IOb } m$ 
and  $v = c \ m$ 
shows  $(v > 0) \wedge (\forall x \ y. ((\exists p. (\text{Ph } p \wedge \text{W } m \ p \ x \wedge \text{W } m \ p \ y))$ 
   $\longleftrightarrow ( \text{space2 } x \ y = (c \ m * c \ m) * \text{time2 } x \ y )))$ 
proof -
have  $v\text{Exists}: (\exists v. \text{cTest } m \ v)$  by ( $\text{metis AxPh assms}(1)$ )
hence  $\text{cTest } m \ (c \ m)$  by ( $\text{metis lemSomeFunc}$ )
thus  $?thesis$  using  $\text{assms}(2)$  by  $\text{simp}$ 
qed

```

```

lemma lemCCone:
  assumes IOb m
    and onCone y (lightcone m x)
  shows  $\exists p. (Ph\ p \wedge W\ m\ p\ x \wedge W\ m\ p\ y)$ 
proof -
  have  $(\exists p. (Ph\ p \wedge W\ m\ p\ x \wedge W\ m\ p\ y))$ 
     $\longleftrightarrow (space2\ x\ y = (c\ m * c\ m) * time2\ x\ y)$ 
    by (simp add: assms(1) lemCProps)
  hence ph-exists:  $(space2\ x\ y = (c\ m * c\ m) * time2\ x\ y) \longrightarrow (\exists p. (Ph\ p \wedge W\ m\ p\ x \wedge W\ m\ p\ y))$ 
    by metis
  define lcmx where lcmx = lightcone m x
  have lcmx-vertex: vertex lcmx = x by (simp add: lcmx-def)
  have lcmx-slope: slope lcmx = c m by (simp add: lcmx-def)
  have onCone y lcmx  $\longrightarrow (space2\ x\ y = (c\ m * c\ m) * time2\ x\ y)$ 
    by (metis lcmx-vertex lcmx-slope onCone.simps)
  hence  $space2\ x\ y = (c\ m * c\ m) * time2\ x\ y$  by (metis lcmx-def assms(2))
  thus ?thesis by (metis ph-exists)
qed

```

```

lemma lemCPos:
  assumes IOb m
  shows  $c\ m > 0$ 
  by (metis assms(1) lemCProps)

```

```

lemma lemCPhoton:
  assumes IOb m
  shows  $\forall x\ y. (\exists p. (Ph\ p \wedge W\ m\ p\ x \wedge W\ m\ p\ y)) \longleftrightarrow (space2\ x\ y = (c\ m * c\ m) * (time2\ x\ y))$ 
    by (metis assms(1) lemCProps)

```

end

```

class AxEv = WorldView +
assumes
  AxEv:  $\llbracket IOb\ m; IOb\ k \rrbracket \implies (\exists y. (\forall b. (W\ m\ b\ x \longleftrightarrow W\ k\ b\ y)))$ 
begin
end

```

```

class AxThExp = WorldView + AxPh +
assumes

```

```

    AxThExp: IOb m  $\implies$  ( $\forall x y .$ 
      ( $\exists k.(IOb k \wedge W m k x \wedge W m k y)$ )  $\longleftrightarrow$  ( $space2 x y < (c m * c m) * time2$ 
 $x y$ )
    ))

```

```

begin
end

```

```

class AxSelf = WorldView +
assumes
  AxSelf: IOb m  $\implies$  ( $W m m x$ )  $\longrightarrow$  ( $onAxisT x$ )
begin
end

```

```

class AxC = WorldView + AxPh +
assumes
  AxC: IOb m  $\implies$   $c m = 1$ 
begin
end

```

```

class AxSym = WorldView +
assumes
  AxSym:  $\llbracket IOb m; IOb k \rrbracket \implies$ 
    ( $W m e x \wedge W m f y \wedge W k e x' \wedge W k f y' \wedge$ 
      $tval x = tval y \wedge tval x' = tval y'$ )
      $\longrightarrow$  ( $space2 x y = space2 x' y'$ )
begin
end

```

```

class AxLines = WorldView +
assumes
  AxLines:  $\llbracket IOb m; IOb k; collinear x p q \rrbracket \implies$ 
     $collinear (wvt k m x) (wvt k m p) (wvt k m q)$ 
begin
end

```

```

class AxPlanes = WorldView +
assumes
  AxPlanes:  $\llbracket \text{IOb } m; \text{IOb } k \rrbracket \implies$ 
    (coplanar e x y z  $\longrightarrow$  coplanar (wvt k m e) (wvt k m x) (wvt k m y) (wvt k m z))
begin
end

```

```

class AxCones = WorldView + AxPh +
assumes
  AxCones:  $\llbracket \text{IOb } m; \text{IOb } k \rrbracket \implies$ 
    (onCone x (lightCone m v)  $\longrightarrow$  onCone (wvt k m x) (lightcone k (wvt k m v)))
begin
end

```

```

class AxTime = WorldView +
assumes
  AxTime:  $\llbracket \text{IOb } m; \text{IOb } k \rrbracket$ 
     $\implies (x \lesssim y \longrightarrow \text{wvt } k \ m \ x \lesssim \text{wvt } k \ m \ y)$ 
begin
end

```

end

```

theory SpecRel
imports Axioms
begin

```

```

class SpecRel = WorldView + AxPh + AxEv + AxSelf + AxSym

```

```

  + AxEuclidean

```

```

  + AxLines + AxPlanes + AxCones

```

```

begin

```

lemma *lemZEG*:
shows $z - e = g - e + (z - g)$
proof –
have $g - e + (z - g) = (g - e + z) - g$ **by** (*rule add-diff-eq*)
also have $(g - e + z) - g = (-e + z)$
by (*metis local.diff-add-cancel*
local.ring-normalization-rules(2)
local.semiring-normalization-rules(24)
local.semiring-normalization-rules(25))
thus *?thesis*
by (*simp add: calculation*)
qed

lemma *noFTLObserver*:
assumes *iobm*: *IOb m*
and *iobk*: *IOb k*
and *mke*: *m sees k at e*
and *mkf*: *m sees k at f*
and *enotf*: $e \neq f$
shows $\text{space2 } e f \leq (c m * c m) * \text{time2 } e f$
proof –

{
assume *converse*: $\text{space2 } e f > (c m * c m) * \text{time2 } e f$

define *eCone* **where** $eCone = mkCone e (c m)$
have *e-on-econe*: *onCone e eCone* **by** (*simp add: eCone-def*)

have *e-is-vertex*: $e = \text{vertex } eCone$ **by** (*simp add: eCone-def*)
have *cm-is-slope*: $c m = \text{slope } eCone$ **by** (*simp add: eCone-def*)
hence *outside*: *outsideCone f eCone*
by (*metis (lifting) e-is-vertex cm-is-slope converse outsideCone.simps*)

have *outsideCone f eCone*
 $\longrightarrow (\exists x. (\text{onCone } x eCone \wedge x \neq \text{vertex } eCone \wedge \text{inPlane } f (\text{tangentPlane } x eCone)))$
by (*rule AxParallelConesE*)

hence *tplane-exists*: $\exists x. (\text{onCone } x eCone \wedge x \neq \text{vertex } eCone \wedge \text{inPlane } f (\text{tangentPlane } x eCone))$

by (*metis outside*)
then obtain g **where** g -props: ($onCone\ g\ eCone \wedge g \neq vertex\ eCone \wedge inPlane\ f\ (tangentPlane\ g\ eCone)$)
by *auto*
have g -on-eCone: $onCone\ g\ eCone$ **by** (*metis g-props*)
have g -not-vertex: $g \neq vertex\ eCone$ **by** (*metis g-props*)

define $tplane$ **where** $tplane = tangentPlane\ g\ eCone$
have e -in-tplane: $inPlane\ e\ tplane$ **by** (*metis AxTangentVertex e-is-vertex tplane-def*)
have f -in-tplane: $inPlane\ f\ tplane$ **by** (*metis g-props tplane-def*)
have g -in-tplane: $inPlane\ g\ tplane$ **by** (*metis lemPlaneContainsBasePoint tplane-def AxTangentBase*)

have ($onCone\ g\ eCone$) \longrightarrow
 $((inPlane\ f\ (tangentPlane\ g\ eCone) \wedge onCone\ f\ eCone)$
 $\longleftrightarrow collinear\ (vertex\ eCone)\ g\ f)$
by (*metis AxConeTangent*)
hence $axconetangent$: $collinear\ e\ g\ f \longrightarrow onCone\ f\ eCone$
by (*metis g-on-eCone e-is-vertex*)
have $\neg(onCone\ f\ eCone)$ **by** (*metis outside lemOutsideNotOnCone*)
hence g -not-collinear: $\neg(collinear\ e\ g\ f)$
by (*metis axconetangent*)

define wte **where** $wte = wvt\ k\ m\ e$
define wtf **where** $wtf = wvt\ k\ m\ f$
define wtg **where** $wtg = wvt\ k\ m\ g$

have $W\ k\ k\ wte$ **by** (*metis wte-def AxWVT mke iobm iobk*)
hence wte -onAxis: $onAxisT\ wte$ **by** (*metis AxSelf iobk*)

have $W\ k\ k\ wtf$ **by** (*metis wtf-def AxWVT mkf iobm iobk*)
hence wtf -onAxis: $onAxisT\ wtf$ **by** (*metis AxSelf iobk*)

have wte -inv: $e = wvt\ m\ k\ wte$ **by** (*metis AxWVTSym iobk iobm wte-def*)
have wtf -inv: $f = wvt\ m\ k\ wtf$ **by** (*metis AxWVTSym iobk iobm wtf-def*)
have wtg -inv: $g = wvt\ m\ k\ wtg$ **by** (*metis AxWVTSym iobk iobm wtg-def*)

have e -not- g : $e \neq g$ **by** (*metis e-is-vertex g-not-vertex*)
have f -not- g : $f \neq g$ **by** (*metis outside lemOutsideNotOnCone g-on-eCone*)

have wt - e -not- f : $wte \neq wtf$ **by** (*metis wte-inv wtf-inv enotf*)
have wt - f -not- g : $wtf \neq wtg$ **by** (*metis wtf-inv wtg-inv f-not-g*)
have wt - g -not- e : $wtg \neq wte$ **by** (*metis wtg-inv wte-inv e-not-g*)

have *if-g-onAxis*: $onAxisT\ wvtg \longrightarrow collinear\ wvte\ wvtg\ wvtf$
by (*metis lemAxisIsLine wvte-onAxis wvtf-onAxis wvt-e-not-f wvt-f-not-g wvt-g-not-e*)

have $collinear\ wvte\ wvtg\ wvtf \longrightarrow collinear\ e\ g\ f$
by (*metis AxLines iobm iobk wvte-inv wvtf-inv wvtg-inv*)
hence $onAxisT\ wvtg \longrightarrow collinear\ e\ g\ f$ **by** (*metis if-g-onAxis*)

hence *wvtg-offAxis*: $\neg (onAxisT\ wvtg)$ **by** (*metis g-not-collinear*)

have $\forall s.(\exists p.(collinear\ wvte\ wvtg\ p \wedge (space2\ p\ wvtf = (s*s)*time2\ p\ wvtf)))$
by (*metis AxSlopedLineInVerticalPlane wvte-onAxis wvtf-onAxis wvtg-offAxis wvt-e-not-f*)
hence *exists-wvtz*: $\exists p.(collinear\ wvte\ wvtg\ p \wedge (space2\ p\ wvtf = (c\ k * c\ k)*time2\ p\ wvtf))$
by *metis*
then obtain *wvtz where*
wvtz-props: $collinear\ wvte\ wvtg\ wvtz \wedge (space2\ wvtz\ wvtf = (c\ k * c\ k)*time2\ wvtz\ wvtf)$ **by** *auto*
hence *wvtz-speed*: $space2\ wvtz\ wvtf = (c\ k * c\ k)*time2\ wvtz\ wvtf$ **by** *metis*

define *z where* $z = wvt\ m\ k\ wvtz$
define *wvtzCone where* $wvtzCone = lightcone\ k\ wvtz$

have *wvtz-is-vertex*: $wvtz = vertex\ wvtzCone$ **by** (*simp add: wvtzCone-def*)
have *ck-is-slope*: $c\ k = slope\ wvtzCone$ **by** (*simp add: wvtzCone-def*)
hence $space2\ (vertex\ wvtzCone)\ wvtf = ((slope\ wvtzCone) * (slope\ wvtzCone))*time2\ (vertex\ wvtzCone)\ wvtf$
by (*metis wvtz-speed wvtz-is-vertex ck-is-slope*)
hence *onCone wvtz wvtzCone* **by** (*metis onCone.simps*)

hence *wvtf-on-wvtzCone*: $onCone\ (wvt\ m\ k\ wvtz)\ (lightcone\ m\ z)$
by (*metis iobm iobk AxCones wvtzCone-def z-def*)

define *zCone where* $zCone = lightcone\ m\ z$
have *z-is-vertex*: $z = vertex\ zCone$ **by** (*simp add: zCone-def*)
have *cm-is-zSlope*: $c\ m = slope\ zCone$ **by** (*simp add: zCone-def*)

have *f-on-zCone*: $onCone\ f\ zCone$ **by** (*metis wvtz-inv wvtf-on-wvtzCone zCone-def*)

hence $\text{space2 } (\text{vertex } z\text{Cone}) f = (\text{slope } z\text{Cone} * \text{slope } z\text{Cone}) * \text{time2 } (\text{vertex } z\text{Cone}) f$
by (*simp add: zCone-def*)
hence $\text{space2 } z f = (c m * c m) * \text{time2 } z f$ **by** (*metis z-is-vertex cm-is-zSlope*)
hence $\text{fz-speed: space2 } f z = (c m * c m) * \text{time2 } f z$ **by** (*metis lemSpace2Sym lemTime2Sym*)

define $f\text{Cone}$ **where** $f\text{Cone} = \text{lightcone } m f$

have $f\text{-is-fVertex: } f = \text{vertex } f\text{Cone}$ **by** (*simp add: fCone-def*)
have $\text{cm-is-fSlope: } c m = \text{slope } f\text{Cone}$ **by** (*simp add: fCone-def*)
hence $\text{space2 } (\text{vertex } f\text{Cone}) z = ((\text{slope } f\text{Cone}) * (\text{slope } f\text{Cone})) * \text{time2 } (\text{vertex } f\text{Cone}) z$
by (*metis fz-speed f-is-fVertex cm-is-fSlope*)
hence $z\text{-on-fCone: onCone } z f\text{Cone}$ **by** (*metis onCone.simps*)

have $\text{collinear } w\text{te } w\text{tg } w\text{tz}$ **by** (*metis wtz-props*)
hence $\text{egz-collinear: collinear } e g z$ **by** (*metis wte-inv wtg-inv z-def AxLines iobm iobk*)
hence $z\text{-geometry: } (\text{inPlane } z (\text{tangentPlane } g e\text{Cone}) \wedge \text{onCone } z e\text{Cone})$
by (*metis AxConeTangent e-is-vertex g-on-eCone*)

have $z\text{-on-eCone: onCone } z e\text{Cone}$ **by** (*metis z-geometry*)
have $z\text{-in-tplane: inPlane } z \text{tplane}$ **by** (*metis z-geometry tplane-def*)

hence $z\text{-not-f: } z \neq f$ **by** (*metis z-on-eCone outside lemOutsideNotOnCone*)
hence $z\text{-not-fVertex: } z \neq \text{vertex } f\text{Cone}$ **by** (*simp add: fCone-def z-not-f*)

{
assume $\text{assm: } z = e$
have $\text{space2 } f e = (c m * c m) * \text{time2 } f e \wedge \text{space2 } f e = \text{space2 } e f \wedge \text{time2 } f e = \text{time2 } e f$
by (*metis lemSpace2Sym lemTime2Sym fz-speed assm*)
hence $\text{space2 } e f = (c m * c m) * \text{time2 } e f$ **by** *metis*
hence False **by** (*metis less-irrefl converse*)
}
from this have $z\text{-not-e: } z \neq e$ **by** *blast*

define lineA **where** $\text{lineA} = \text{lineJoining } e z$
define lineB **where** $\text{lineB} = \text{lineJoining } f z$

{

```

    assume assm: direction lineA = vecZero
    have lemnullline: (direction lineA = vecZero  $\wedge$  inLine z lineA)  $\longrightarrow$  z = basepoint
lineA
      by (metis lemNullLine)
    have inLine z lineA by (metis lineA-def lemLineContainsEndpoint)
    hence z-is-bp: z = basepoint lineA by (metis lemnullline assm)
    have basepoint lineA = e by (simp add: lineA-def)
    hence False by (metis z-is-bp z-not-e)
  }
from this have ez-not-null: direction lineA  $\neq$  vecZero by blast

{
  assume assm: direction lineB = vecZero
  have lemnullline: (direction lineB = vecZero  $\wedge$  inLine z lineB)  $\longrightarrow$  z = basepoint
lineB
    by (metis lemNullLine)
  have inLine z lineB by (metis lineB-def lemLineContainsEndpoint)
  hence z-is-bp: z = basepoint lineB by (metis lemnullline assm)
  have basepoint lineB = f by (simp add: lineB-def)
  hence False by (metis z-is-bp z-not-f)
}
from this have fz-not-null: direction lineB  $\neq$  vecZero by blast

{
  have samePlane tplane (tangentPlane z fCone
     $\wedge$  (lineJoining e g  $\parallel$  lineJoining f z))
  by (metis AxParallelCones tplane-def
    g-on-eCone g-not-vertex z-on-fCone z-not-fVertex z-in-tplane
    e-is-vertex f-is-fVertex)

  hence eg-par-fz: (lineJoining e g  $\parallel$  lineJoining f z) by metis
  {
    assume case1: direction (lineJoining e g) = vecZero
    have direction (lineJoining e g) = from e to g by simp
    hence from e to g = vecZero by (metis case1)
    hence e = g by (simp)
    hence False by (metis e-not-g)
  }
  from this have eg-not-null:  $\neg$ (direction (lineJoining e g) = vecZero) by blast
  then obtain a where a-props: a  $\neq$  0  $\wedge$  direction (lineJoining f z) = a**direction
(lineJoining e g)
    by (metis fz-not-null eg-not-null eg-par-fz parallel.simps lineB-def)
  hence f-to-z: from f to z = a*(from e to g) by simp
  have a-nonzero: a  $\neq$  0 by (metis a-props)

  have eg-dir: from e to g = direction (lineJoining e g) by simp
  have gz-dir: from g to z = direction (lineJoining g z) by simp
  have egz: z = g  $\rightsquigarrow$  (from g to z) by (metis lemLineEndpoint)
  hence collinear e g (g  $\rightsquigarrow$  (from g to z)) by (metis egz-collinear)

```

then obtain b where $e\text{-to-}g$: $\text{from } e \text{ to } g = (-b)(\text{from } g \text{ to } z)$**
by (metis lemDirectionCollinear)

{
assume asm : $-b = 0$
have $\text{from } e \text{ to } g = (-b)(\text{from } g \text{ to } z)$ **by (metis e-to-g)****
hence $\text{from } e \text{ to } g = \text{vecZero}$ **by (simp add: asm)**
hence $\text{direction } (\text{lineJoining } e \ g) = \text{vecZero}$ **by (simp)**
hence False **by (metis eg-not-null lineA-def)**
}

from this have $b\text{-nonzero}$: $-b \neq 0$ **by blast**

define $binv$ where $binv = \text{inverse } (-b)$
define $factor$ where $factor = 1 + binv$
have $binv\text{-nonzero}$: $binv \neq 0$ **by (metis b-nonzero add.comm-neutral binv-def
nonzero-imp-inverse-nonzero right-minus)**

have $\text{from } e \text{ to } g = (-b)(\text{from } g \text{ to } z)$ **by (metis e-to-g)****
hence $g\text{-to-}z$: $(\text{from } g \text{ to } z) = binv(\text{from } e \text{ to } g)$**
****by (metis b-nonzero lemScaleInverse binv-def)****

have $\text{from } e \text{ to } z = \text{from } e \text{ to } g \oplus \text{from } g \text{ to } z$
****by (simp add: lemZEG)****

hence $\text{from } e \text{ to } z = (\text{from } e \text{ to } g) \oplus binv(\text{from } e \text{ to } g)$ **by (metis g-to-z)****
hence $e\text{-to-}z$: $\text{from } e \text{ to } z = factor(\text{from } e \text{ to } g)$ **by (metis lemAddOverScale**
lemScale1 factor-def)**

have $ez\text{-dir}$: $\text{direction } (\text{lineJoining } e \ z) = \text{from } e \text{ to } z$ **by simp**
have $eg\text{-dir}$: $\text{direction } (\text{lineJoining } e \ g) = \text{from } e \text{ to } g$ **by simp**

{
assume asm : $factor = 0$
have $\text{from } e \text{ to } z = factor(\text{from } e \text{ to } g)$ **by (metis e-to-z)****
hence $\text{from } e \text{ to } z = \text{vecZero}$ **by (simp add: asm)**
hence $\text{direction } (\text{lineJoining } e \ z) = \text{vecZero}$ **by (simp)**
hence False **by (metis ez-not-null lineA-def)**
}

from this have $factor\text{-nonzero}$: $factor \neq 0$ **by blast**

have $\text{direction } (\text{lineJoining } e \ z) = factor(\text{direction } (\text{lineJoining } e \ g))$**
****by (metis e-to-z ez-dir eg-dir)****
hence $(\text{lineJoining } e \ g) \parallel (\text{lineJoining } e \ z)$ **by (metis parallel.simps fac-
tor-nonzero)**

hence $(\text{lineJoining } e \ z) \parallel (\text{lineJoining } e \ g)$ **by (metis lemParallelSym)**

hence $(\text{lineJoining } e \ z) \parallel (\text{lineJoining } f \ z)$ **by (metis lemParallelTrans eg-par-fz)**

```

eg-not-null)
}
from this have A-par-B: lineA || lineB by (metis lineA-def lineB-def)

have e-in-lineA: inLine e lineA by (metis lineA-def lemLineContainsBasepoint)

{
have basic:  $\forall a b. (((-a)*b)*((-a)*b) = (a*a)*(b*b))$ 
by (metis equation-minus-iff minus-mult-commute minus-mult-right
semiring-normalization-rules(17) semiring-normalization-rules(19))

assume assm: inLine e lineB
hence coll: collinear e f (f  $\rightsquigarrow$  direction lineB) by (simp add: lineB-def)
then obtain  $\beta$  where props: from e to f =  $(-\beta)**(\text{direction lineB})$ 
by (metis lemDirectionCollinear)

hence tval f - tval e =  $(-\beta)*(tval z - tval f) \wedge xval f - xval e = (-\beta)*(xval z - xval f)$ 
 $\wedge yval f - yval e = (-\beta)*(yval z - yval f) \wedge zval f - zval e = (-\beta)*(zval z - zval f)$ 
by (simp add: lineB-def)
hence speeds: time2 f e =  $(\beta*\beta)*time2 z f \wedge space2 f e = (\beta*\beta)*space2 z f$ 
apply (simp add: basic) apply auto
apply (metis semiring-normalization-rules(18) semiring-normalization-rules(19))
by (metis semiring-normalization-rules(18) semiring-normalization-rules(19))

semiring-normalization-rules(34))

have space2 f z =  $(c m * c m)*time2 f z$  by (metis fz-speed)
hence space2 z f =  $(c m * c m)*time2 z f$  by (metis lemSpace2Sym lem-
Time2Sym)
hence space2 f e =  $((\beta*\beta)*(c m * c m))*time2 z f$  by (metis speeds mult.assoc)
hence space2 f e =  $(c m * c m)*(\beta*\beta)*time2 z f$  by (metis mult.assoc
mult.commute)
hence space2 f e =  $(c m * c m)*time2 f e$  by (metis mult.assoc speeds)
hence space2 e f =  $(c m * c m)*time2 e f$  by (metis lemSpace2Sym lem-
Time2Sym)
hence False by (metis less-irrefl converse)
}
from this have e-not-in-lineB:  $\neg(\text{inLine } e \text{ lineB})$  by blast

have inLine z lineA  $\wedge$  inLine z lineB by (metis lemLineContainsEndpoint lineA-def
lineB-def)
hence A-meets-B: meets lineA lineB by auto

hence False by (metis A-par-B ez-not-null fz-not-null e-in-lineA e-not-in-lineB

```

```

lemParallelNotMeet)
}
from this have  $\neg (\text{space2 } e f > (c\ m * c\ m) * \text{time2 } e f)$  by blast

thus ?thesis by simp
qed

end

end

```

References

- [1] M. Stannett and I. Németi. Using Isabelle/HOL to verify first-order relativity theory. *Journal of Automated Reasoning*, 52(4):361–378, 2014.