

Interval Temporal Logic on Natural Numbers

David Trachtenherz

October 13, 2025

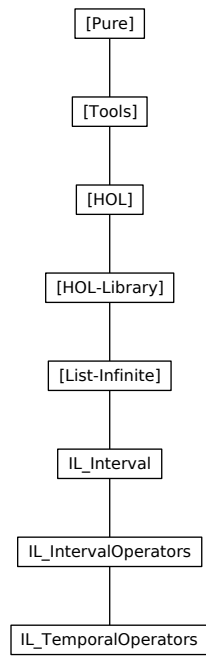
Abstract

We introduce a theory of temporal logic operators using sets of natural numbers as time domain, formalized in a shallow embedding manner. The theory comprises special natural intervals (theory `IL_Interval`: open and closed intervals, continuous and modulo intervals, interval traversing results), operators for shifting intervals to left/right on the number axis as well as expanding/contracting intervals by constant factors (theory `IL_IntervalOperators.thy`), and ultimately definitions and results for unary and binary temporal operators on arbitrary natural sets (theory `IL_TemporalOperators`).

Contents

1	Intervals and operations for temporal logic declarations	2
1.1	Time intervals – definitions and basic lemmata	2
1.1.1	Definitions	2
1.1.2	Membership in an interval	3
1.1.3	Interval conversions	5
1.1.4	Finiteness and emptiness of intervals	6
1.1.5	<i>Min</i> and <i>Max</i> element of an interval	7
1.2	Adding and subtracting constants to interval elements	8
1.3	Relations between intervals	11
1.3.1	Auxiliary lemmata	11
1.3.2	Subset relation between intervals	12
1.3.3	Equality of intervals	18
1.3.4	Inequality of intervals	18
1.4	Union and intersection of intervals	19
1.5	Cutting intervals	23
1.6	Cardinality of intervals	30
1.7	Functions <i>inext</i> and <i>iprev</i> with intervals	32
1.7.1	Mirroring of intervals	35
1.7.2	Functions <i>inext-nth</i> and <i>iprev-nth</i> on intervals	36
1.8	Induction with intervals	38

2	Arithmetic operators on natural intervals	39
2.1	Arithmetic operations with intervals	39
2.1.1	Addition of and multiplication by constants	39
2.1.2	Some conversions between intervals using constant ad- dition and multiplication	43
2.1.3	Subtraction of constants	44
2.1.4	Subtraction of intervals from constants	49
2.1.5	Division of intervals by constants	52
2.2	Interval cut operators with arithmetic interval operators . . .	59
2.3	<i>inext</i> and <i>iprev</i> with interval operators	61
2.4	Cardinality of intervals with interval operators	64
2.5	Results about sets of intervals	67
2.5.1	Set of intervals without and with empty interval . . .	67
2.5.2	Interval sets are closed under cutting	72
2.5.3	Interval sets are closed under addition and multiplication	73
2.5.4	Interval sets are closed with certain conditions under subtraction	73
2.5.5	Interval sets are not closed under division	74
2.5.6	Sets of intervals closed under division	74
3	Temporal logic operators on natural intervals	78
3.1	Basic definitions	79
3.2	Basic lemmata for temporal operators	82
3.2.1	Intro/elim rules	82
3.2.2	Rewrite rules for trivial simplification	83
3.2.3	Empty sets and singletons	89
3.2.4	Conversions between temporal operators	90
3.2.5	Some implication results	93
3.2.6	Congruence rules for temporal operators' predicates .	94
3.2.7	Temporal operators with set unions/intersections and subsets	95
3.3	Further results for temporal operators	96
3.4	Temporal operators and arithmetic interval operators	98



1 Intervals and operations for temporal logic declarations

```

theory IL-Interval
imports
  List-Infinite.InfiniteSet2
  List-Infinite.SetIntervalStep
begin

```

1.1 Time intervals – definitions and basic lemmata

1.1.1 Definitions

```

type-synonym Time = nat

```

```

type-synonym iT = Time set

```

Infinite interval starting at some natural n .

definition

```

iFROM :: Time  $\Rightarrow$  iT ( $\langle [-..] \rangle$ )

```

where

```

 $[-..] \equiv \{n.. \}$ 

```

Finite interval starting at 0 and ending at some natural n .

definition

```

iTILL :: Time  $\Rightarrow$  iT ( $\langle [..n] \rangle$ )

```

where

```

 $[..n] \equiv \{..n\}$ 

```

Finite bounded interval containing the naturals between n and $n + d$. d denotes the difference between left and right interval bound. The number of elements is $d + 1$ so that an empty interval cannot be defined.

definition

```

iIN :: Time  $\Rightarrow$  nat  $\Rightarrow$  iT ( $\langle [-.., -] \rangle$ )

```

where

```

 $[n.., d] \equiv \{n..n+d\}$ 

```

Infinite modulo interval containing all naturals having the same division remainder modulo m as r , and beginning at n .

definition

```

iMOD :: Time  $\Rightarrow$  nat  $\Rightarrow$  iT ( $\langle [ -, mod - ] \rangle$ )

```

where

```

 $[r, mod m] \equiv \{ x. x \bmod m = r \bmod m \wedge r \leq x \}$ 

```

Finite bounded modulo interval containing all naturals having the same division remainder modulo m as r , beginning at n , and ending after c cycles

at $r + m * c$. The number of elements is $c + 1$ so that an empty interval cannot be defined.

definition

$iMODb :: Time \Rightarrow nat \Rightarrow nat \Rightarrow iT \ (\langle [-, mod -, -] \rangle)$

where

$[r, mod\ m, c] \equiv \{ x. x\ mod\ m = r\ mod\ m \wedge r \leq x \wedge x \leq r + m * c \}$

1.1.2 Membership in an interval

lemmas $iT-defs = iFROM-def\ iTILL-def\ iIN-def\ iMOD-def\ iMODb-def$

lemma $iFROM-iff: x \in [n..] = (n \leq x)$

$\langle proof \rangle$

lemma $iTILL-iff: x \in [...n] = (x \leq n)$

$\langle proof \rangle$

lemma $iIN-iff: x \in [n...,d] = (n \leq x \wedge x \leq n + d)$

$\langle proof \rangle$

lemma $iMOD-iff: x \in [r, mod\ m] = (x\ mod\ m = r\ mod\ m \wedge r \leq x)$

$\langle proof \rangle$

lemma $iMODb-iff: x \in [r, mod\ m, c] =$

$(x\ mod\ m = r\ mod\ m \wedge r \leq x \wedge x \leq r + m * c)$

$\langle proof \rangle$

lemma $iFROM-D: x \in [n..] \implies (n \leq x)$

$\langle proof \rangle$

lemma $iTILL-D: x \in [...n] \implies (x \leq n)$

$\langle proof \rangle$

corollary $iIN-geD: x \in [n...,d] \implies n \leq x$

$\langle proof \rangle$

corollary $iIN-leD: x \in [n...,d] \implies x \leq n + d$

$\langle proof \rangle$

corollary $iMOD-modD: x \in [r, mod\ m] \implies x\ mod\ m = r\ mod\ m$

$\langle proof \rangle$

corollary $iMOD-geD: x \in [r, mod\ m] \implies r \leq x$

$\langle proof \rangle$

corollary $iMODb-modD: x \in [r, mod\ m, c] \implies x\ mod\ m = r\ mod\ m$

$\langle proof \rangle$

corollary $iMODb-geD: x \in [r, mod\ m, c] \implies r \leq x$

$\langle proof \rangle$

corollary $iMODb-leD: x \in [r, mod\ m, c] \implies x \leq r + m * c$

$\langle proof \rangle$

lemmas $iT-iff = iFROM-iff\ iTILL-iff\ iIN-iff\ iMOD-iff\ iMODb-iff$

lemmas $iT-drule =$

$iFROM-D$

$iTILL-D$

$iIN-geD\ iIN-leD$

$iMOD-modD\ iMOD-geD$

iMODb-modD iMODb-geD iMODb-leD

lemma

iFROM-I [intro]: $n \leq x \implies x \in [n..]$ **and**
iTILL-I [intro]: $x \leq n \implies x \in [...n]$ **and**
iIN-I [intro]: $n \leq x \implies x \leq n + d \implies x \in [n..,d]$ **and**
iMOD-I [intro]: $x \bmod m = r \bmod m \implies r \leq x \implies x \in [r, \bmod m]$ **and**
iMODb-I [intro]: $x \bmod m = r \bmod m \implies r \leq x \implies x \leq r + m * c \implies x \in [r, \bmod m, c]$
 <proof>

lemma

iFROM-E [elim]: $x \in [n..] \implies (n \leq x \implies P) \implies P$ **and**
iTILL-E [elim]: $x \in [...n] \implies (x \leq n \implies P) \implies P$ **and**
iIN-E [elim]: $x \in [n..,d] \implies (n \leq x \implies x \leq n + d \implies P) \implies P$ **and**
iMOD-E [elim]: $x \in [r, \bmod m] \implies (x \bmod m = r \bmod m \implies r \leq x \implies P) \implies P$ **and**
iMODb-E [elim]: $x \in [r, \bmod m, c] \implies (x \bmod m = r \bmod m \implies r \leq x \implies x \leq r + m * c \implies P) \implies P$
 <proof>

lemma *iIN-Suc-insert-conv*:

$\text{insert } (\text{Suc } (n + d)) [n..,d] = [n..,\text{Suc } d]$
 <proof>

lemma *iTILL-Suc-insert-conv*: $\text{insert } (\text{Suc } n) [...n] = [...\text{Suc } n]$

<proof>

lemma *iMODb-Suc-insert-conv*:

$\text{insert } (r + m * \text{Suc } c) [r, \bmod m, c] = [r, \bmod m, \text{Suc } c]$
 <proof>

lemma *iFROM-pred-insert-conv*: $\text{insert } (n - \text{Suc } 0) [n..] = [n - \text{Suc } 0..]$

<proof>

lemma *iIN-pred-insert-conv*:

$0 < n \implies \text{insert } (n - \text{Suc } 0) [n..,d] = [n - \text{Suc } 0..,\text{Suc } d]$
 <proof>

lemma *iMOD-pred-insert-conv*:

$m \leq r \implies \text{insert } (r - m) [r, \bmod m] = [r - m, \bmod m]$
 <proof>

lemma *iMODb-pred-insert-conv*:

$m \leq r \implies \text{insert } (r - m) [r, \bmod m, c] = [r - m, \bmod m, \text{Suc } c]$

$\langle \text{proof} \rangle$

lemma *iFROM-Suc-pred-insert-conv*: $\text{insert } n \text{ [Suc } n \dots] = [n \dots]$

$\langle \text{proof} \rangle$

lemma *iIN-Suc-pred-insert-conv*: $\text{insert } n \text{ [Suc } n \dots, d] = [n \dots, \text{Suc } d]$

$\langle \text{proof} \rangle$

lemma *iMOD-Suc-pred-insert-conv*: $\text{insert } r \text{ [} r + m, \text{ mod } m] = [r, \text{ mod } m]$

$\langle \text{proof} \rangle$

lemma *iMODb-Suc-pred-insert-conv*: $\text{insert } r \text{ [} r + m, \text{ mod } m, c] = [r, \text{ mod } m, \text{Suc } c]$

$\langle \text{proof} \rangle$

lemmas *iT-Suc-insert* =

iIN-Suc-insert-conv

iTILL-Suc-insert-conv

iMODb-Suc-insert-conv

lemmas *iT-pred-insert* =

iFROM-pred-insert-conv

iIN-pred-insert-conv

iMOD-pred-insert-conv

iMODb-pred-insert-conv

lemmas *iT-Suc-pred-insert* =

iFROM-Suc-pred-insert-conv

iIN-Suc-pred-insert-conv

iMOD-Suc-pred-insert-conv

iMODb-Suc-pred-insert-conv

lemma *iMOD-mem-diff*: $\llbracket a \in [r, \text{ mod } m]; b \in [r, \text{ mod } m] \rrbracket \implies (a - b) \text{ mod } m = 0$

$\langle \text{proof} \rangle$

lemma *iMODb-mem-diff*: $\llbracket a \in [r, \text{ mod } m, c]; b \in [r, \text{ mod } m, c] \rrbracket \implies (a - b) \text{ mod } m = 0$

$\langle \text{proof} \rangle$

1.1.3 Interval conversions

lemma *iIN-0-iTILL-conv*: $[0 \dots, n] = [\dots n]$

$\langle \text{proof} \rangle$

lemma *iIN-iTILL-iTILL-conv*: $0 < n \implies [n \dots, d] = [\dots n + d] - [\dots n - \text{Suc } 0]$

$\langle \text{proof} \rangle$

lemma *iIN-iFROM-iTILL-conv*: $[n \dots, d] = [n \dots] \cap [\dots n + d]$

$\langle \text{proof} \rangle$

lemma *iMODb-iMOD-iTILL-conv*: $[r, \text{ mod } m, c] = [r, \text{ mod } m] \cap [\dots r + m * c]$

$\langle \text{proof} \rangle$

lemma *iMODb-iMOD-iIN-conv*: $[r, \text{ mod } m, c] = [r, \text{ mod } m] \cap [r \dots, m * c]$

$\langle \text{proof} \rangle$

lemma *iFROM-iTILL-iIN-conv*: $n \leq n' \implies [n \dots] \cap [\dots n'] = [n \dots, n' - n]$

$\langle \text{proof} \rangle$

lemma *iMOD-iTILL-iMODb-conv*:

$$r \leq n \implies [r, \text{mod } m] \cap [\dots n] = [r, \text{mod } m, (n - r) \text{ div } m]$$

<proof>

lemma *iMOD-iIN-iMODb-conv*:

$$[r, \text{mod } m] \cap [r\dots, d] = [r, \text{mod } m, d \text{ div } m]$$

<proof>

lemma *iFROM-0*: $[0\dots] = \text{UNIV}$

<proof>

lemma *iTILL-0*: $[\dots 0] = \{0\}$

<proof>

lemma *iIN-0*: $[n\dots, 0] = \{n\}$

<proof>

lemma *iMOD-0*: $[r, \text{mod } 0] = [r\dots, 0]$

<proof>

lemma *iMODb-mod-0*: $[r, \text{mod } 0, c] = [r\dots, 0]$

<proof>

lemma *iMODb-0*: $[r, \text{mod } m, 0] = [r\dots, 0]$

<proof>

lemmas *iT-0* =

iFROM-0

iTILL-0

iIN-0

iMOD-0

iMODb-mod-0

iMODb-0

lemma *iMOD-1*: $[r, \text{mod } \text{Suc } 0] = [r\dots]$

<proof>

lemma *iMODb-mod-1*: $[r, \text{mod } \text{Suc } 0, c] = [r\dots, c]$

<proof>

1.1.4 Finiteness and emptiness of intervals

lemma

iFROM-not-empty: $[n\dots] \neq \{\}$ **and**

iTILL-not-empty: $[\dots n] \neq \{\}$ **and**

iIN-not-empty: $[n\dots, d] \neq \{\}$ **and**

iMOD-not-empty: $[r, \text{mod } m] \neq \{\}$ **and**

iMODb-not-empty: $[r, \text{mod } m, c] \neq \{\}$

<proof>

lemmas *iT-not-empty* =
iFROM-not-empty
iTILL-not-empty
iIN-not-empty
iMOD-not-empty
iMODb-not-empty

lemma
iTILL-finite: *finite* $[\dots n]$ **and**
iIN-finite: *finite* $[n\dots, d]$ **and**
iMODb-finite: *finite* $[r, \text{mod } m, c]$ **and**
iMOD-0-finite: *finite* $[r, \text{mod } 0]$
 $\langle \text{proof} \rangle$

lemma *iFROM-infinite*: *infinite* $[n\dots]$
 $\langle \text{proof} \rangle$

lemma *iMOD-infinite*: $0 < m \implies \text{infinite } [r, \text{mod } m]$
 $\langle \text{proof} \rangle$

lemmas *iT-finite* =
iTILL-finite
iIN-finite
iMODb-finite *iMOD-0-finite*

lemmas *iT-infinite* =
iFROM-infinite
iMOD-infinite

1.1.5 *Min* and *Max* element of an interval

lemma
iTILL-Min: *iMin* $[\dots n] = 0$ **and**
iFROM-Min: *iMin* $[n\dots] = n$ **and**
iIN-Min: *iMin* $[n\dots, d] = n$ **and**
iMOD-Min: *iMin* $[r, \text{mod } m] = r$ **and**
iMODb-Min: *iMin* $[r, \text{mod } m, c] = r$
 $\langle \text{proof} \rangle$

lemmas *iT-Min* =
iIN-Min
iTILL-Min
iFROM-Min
iMOD-Min
iMODb-Min

lemma
iTILL-Max: *Max* $[\dots n] = n$ **and**

iIN-Max: $\text{Max } [n.., d] = n + d$ **and**
iMODb-Max: $\text{Max } [r, \text{mod } m, c] = r + m * c$ **and**
iMOD-0-Max: $\text{Max } [r, \text{mod } 0] = r$
 <proof>

lemmas *iT-Max* =
iTILL-Max
iIN-Max
iMODb-Max
iMOD-0-Max

lemma
iTILL-iMax: $i\text{Max } [..n] = \text{enat } n$ **and**
iIN-iMax: $i\text{Max } [n.., d] = \text{enat } (n + d)$ **and**
iMODb-iMax: $i\text{Max } [r, \text{mod } m, c] = \text{enat } (r + m * c)$ **and**
iMOD-0-iMax: $i\text{Max } [r, \text{mod } 0] = \text{enat } r$ **and**
iFROM-iMax: $i\text{Max } [n..] = \infty$ **and**
iMOD-iMax: $0 < m \implies i\text{Max } [r, \text{mod } m] = \infty$
 <proof>

lemmas *iT-iMax* =
iTILL-iMax
iIN-iMax
iMODb-iMax
iMOD-0-iMax
iFROM-iMax
iMOD-iMax

1.2 Adding and subtracting constants to interval elements

lemma
iFROM-plus: $x \in [n..] \implies x + k \in [n..]$ **and**
iFROM-Suc: $x \in [n..] \implies \text{Suc } x \in [n..]$ **and**
iFROM-minus: $\llbracket x \in [n..]; k \leq x - n \rrbracket \implies x - k \in [n..]$ **and**
iFROM-pred: $n < x \implies x - \text{Suc } 0 \in [n..]$
 <proof>

lemma
iTILL-plus: $\llbracket x \in [..n]; k \leq n - x \rrbracket \implies x + k \in [..n]$ **and**
iTILL-Suc: $x < n \implies \text{Suc } x \in [..n]$ **and**
iTILL-minus: $x \in [..n] \implies x - k \in [..n]$ **and**
iTILL-pred: $x \in [..n] \implies x - \text{Suc } 0 \in [..n]$
 <proof>

lemma *iIN-plus*: $\llbracket x \in [n.., d]; k \leq n + d - x \rrbracket \implies x + k \in [n.., d]$
 <proof>

lemma *iIN-Suc*: $\llbracket x \in [n.., d]; x < n + d \rrbracket \implies \text{Suc } x \in [n.., d]$
 <proof>

lemma *iIN-minus*: $\llbracket x \in [n..d]; k \leq x - n \rrbracket \implies x - k \in [n..d]$
 $\langle proof \rangle$

lemma *iIN-pred*: $\llbracket x \in [n..d]; n < x \rrbracket \implies x - Suc\ 0 \in [n..d]$
 $\langle proof \rangle$

lemma *iMOD-plus-divisor-mult*: $x \in [r, mod\ m] \implies x + k * m \in [r, mod\ m]$
 $\langle proof \rangle$

corollary *iMOD-plus-divisor*: $x \in [r, mod\ m] \implies x + m \in [r, mod\ m]$
 $\langle proof \rangle$

lemma *iMOD-minus-divisor-mult*:
 $\llbracket x \in [r, mod\ m]; k * m \leq x - r \rrbracket \implies x - k * m \in [r, mod\ m]$
 $\langle proof \rangle$

corollary *iMOD-minus-divisor-mult2*:
 $\llbracket x \in [r, mod\ m]; k \leq (x - r) \div m \rrbracket \implies x - k * m \in [r, mod\ m]$
 $\langle proof \rangle$

corollary *iMOD-minus-divisor*:
 $\llbracket x \in [r, mod\ m]; m + r \leq x \rrbracket \implies x - m \in [r, mod\ m]$
 $\langle proof \rangle$

lemma *iMOD-plus*:
 $x \in [r, mod\ m] \implies (x + k \in [r, mod\ m]) = (k \bmod m = 0)$
 $\langle proof \rangle$

corollary *iMOD-Suc*:
 $x \in [r, mod\ m] \implies (Suc\ x \in [r, mod\ m]) = (m = Suc\ 0)$
 $\langle proof \rangle$

lemma *iMOD-minus*:
 $\llbracket x \in [r, mod\ m]; k \leq x - r \rrbracket \implies (x - k \in [r, mod\ m]) = (k \bmod m = 0)$
 $\langle proof \rangle$

corollary *iMOD-pred*:
 $\llbracket x \in [r, mod\ m]; r < x \rrbracket \implies (x - Suc\ 0 \in [r, mod\ m]) = (m = Suc\ 0)$
 $\langle proof \rangle$

lemma *iMODb-plus-divisor-mult*:
 $\llbracket x \in [r, mod\ m, c]; k * m \leq r + m * c - x \rrbracket \implies x + k * m \in [r, mod\ m, c]$
 $\langle proof \rangle$

lemma *iMODb-plus-divisor-mult2*:
 $\llbracket x \in [r, mod\ m, c]; k \leq c - (x - r) \div m \rrbracket \implies$
 $x + k * m \in [r, mod\ m, c]$
 $\langle proof \rangle$

lemma *iMODb-plus-divisor*:

$\llbracket x \in [r, \text{mod } m, c]; x < r + m * c \rrbracket \implies x + m \in [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

lemma *iMODb-minus-divisor-mult*:

$\llbracket x \in [r, \text{mod } m, c]; r + k * m \leq x \rrbracket \implies x - k * m \in [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

lemma *iMODb-plus*:

$\llbracket x \in [r, \text{mod } m, c]; k \leq r + m * c - x \rrbracket \implies$
 $(x + k \in [r, \text{mod } m, c]) = (k \text{ mod } m = 0)$
 $\langle \text{proof} \rangle$

corollary *iMODb-Suc*:

$\llbracket x \in [r, \text{mod } m, c]; x < r + m * c \rrbracket \implies$
 $(\text{Suc } x \in [r, \text{mod } m, c]) = (m = \text{Suc } 0)$
 $\langle \text{proof} \rangle$

lemma *iMODb-minus*:

$\llbracket x \in [r, \text{mod } m, c]; k \leq x - r \rrbracket \implies$
 $(x - k \in [r, \text{mod } m, c]) = (k \text{ mod } m = 0)$
 $\langle \text{proof} \rangle$

corollary *iMODb-pred*:

$\llbracket x \in [r, \text{mod } m, c]; r < x \rrbracket \implies$
 $(x - \text{Suc } 0 \in [r, \text{mod } m, c]) = (m = \text{Suc } 0)$
 $\langle \text{proof} \rangle$

lemmas *iFROM-plus-minus* =

iFROM-plus
iFROM-Suc
iFROM-minus
iFROM-pred

lemmas *iTILL-plus-minus* =

iTILL-plus
iTILL-Suc
iTILL-minus
iTILL-pred

lemmas *iIN-plus-minus* =

iIN-plus
iIN-Suc
iTILL-minus
iIN-pred

lemmas *iMOD-plus-minus-divisor* =

iMOD-plus-divisor-mult
iMOD-plus-divisor
iMOD-minus-divisor-mult

iMOD-minus-divisor-mult2
iMOD-minus-divisor

lemmas *iMOD-plus-minus* =
iMOD-plus
iMOD-Suc
iMOD-minus
iMOD-pred

lemmas *iMODb-plus-minus-divisor* =
iMODb-plus-divisor-mult
iMODb-plus-divisor-mult2
iMODb-plus-divisor
iMODb-minus-divisor-mult

lemmas *iMODb-plus-minus* =
iMODb-plus
iMODb-Suc
iMODb-minus
iMODb-pred

lemmas *iT-plus-minus* =
iFROM-plus-minus
iTILL-plus-minus
iIN-plus-minus
iMOD-plus-minus-divisor
iMOD-plus-minus
iMODb-plus-minus-divisor
iMODb-plus-minus

1.3 Relations between intervals

1.3.1 Auxiliary lemmata

lemma *Suc-in-imp-not-subset-iMOD*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies \neg S \subseteq [r, \text{mod } m]$
 $\langle \text{proof} \rangle$

corollary *Suc-in-imp-neq-iMOD*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies S \neq [r, \text{mod } m]$
 $\langle \text{proof} \rangle$

lemma *Suc-in-imp-not-subset-iMODb*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies \neg S \subseteq [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

corollary *Suc-in-imp-neq-iMODb*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies S \neq [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

1.3.2 Subset relation between intervals

lemma

iIN-iFROM-subset-same: $[n \dots, d] \subseteq [n \dots]$ **and**
iIN-iTILL-subset-same: $[n \dots, d] \subseteq [\dots n + d]$ **and**
iMOD-iFROM-subset-same: $[r, \text{mod } m] \subseteq [r \dots]$ **and**
iMODb-iTILL-subset-same: $[r, \text{mod } m, c] \subseteq [\dots r + m * c]$ **and**
iMODb-iIN-subset-same: $[r, \text{mod } m, c] \subseteq [r \dots, m * c]$ **and**
iMODb-iMOD-subset-same: $[r, \text{mod } m, c] \subseteq [r, \text{mod } m]$

$\langle \text{proof} \rangle$

lemmas *iT-subset-same* =

iIN-iFROM-subset-same
iIN-iTILL-subset-same
iMOD-iFROM-subset-same
iMODb-iTILL-subset-same
iMODb-iIN-subset-same
iMODb-iTILL-subset-same
iMODb-iMOD-subset-same

lemma *iMODb-imp-iMOD*: $x \in [r, \text{mod } m, c] \implies x \in [r, \text{mod } m]$

$\langle \text{proof} \rangle$

lemma *iMOD-imp-iMODb*:

$\llbracket x \in [r, \text{mod } m]; x \leq r + m * c \rrbracket \implies x \in [r, \text{mod } m, c]$

$\langle \text{proof} \rangle$

lemma *iMOD-singleton-subset-conv*: $([r, \text{mod } m] \subseteq \{a\}) = (r = a \wedge m = 0)$

$\langle \text{proof} \rangle$

lemma *iMOD-singleton-eq-conv*: $([r, \text{mod } m] = \{a\}) = (r = a \wedge m = 0)$

$\langle \text{proof} \rangle$

lemma *iMODb-singleton-subset-conv*:

$([r, \text{mod } m, c] \subseteq \{a\}) = (r = a \wedge (m = 0 \vee c = 0))$

$\langle \text{proof} \rangle$

lemma *iMODb-singleton-eq-conv*:

$([r, \text{mod } m, c] = \{a\}) = (r = a \wedge (m = 0 \vee c = 0))$

$\langle \text{proof} \rangle$

lemma *iMODb-subset-imp-divisor-mod-0*:

$\llbracket 0 < c'; [r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c] \rrbracket \implies m' \text{ mod } m = 0$

$\langle \text{proof} \rangle$

lemma *iMOD-subset-imp-divisor-mod-0*:

$[r', \text{mod } m'] \subseteq [r, \text{mod } m] \implies m' \text{ mod } m = 0$

$\langle \text{proof} \rangle$

lemma *iMOD-subset-imp-iMODb-subset*:

$\llbracket [r', \text{mod } m'] \subseteq [r, \text{mod } m]; r' + m' * c' \leq r + m * c \rrbracket \implies$
 $[r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]$

$\langle proof \rangle$

lemma *iMODb-subset-imp-iMOD-subset*:

$$\llbracket [r', \text{mod } m', c] \subseteq [r, \text{mod } m, c]; 0 < c' \rrbracket \implies \\ [r', \text{mod } m'] \subseteq [r, \text{mod } m]$$

$\langle proof \rangle$

lemma *iMODb-0-iMOD-subset-conv*:

$$([r', \text{mod } m', 0] \subseteq [r, \text{mod } m]) = \\ (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r')$$

$\langle proof \rangle$

lemma *iFROM-subset-conv*: $([n' \dots] \subseteq [n \dots]) = (n \leq n')$

$\langle proof \rangle$

lemma *iFROM-iMOD-subset-conv*: $([n' \dots] \subseteq [r, \text{mod } m]) = (r \leq n' \wedge m = \text{Suc } 0)$

$\langle proof \rangle$

lemma *iIN-subset-conv*: $([n' \dots, d'] \subseteq [n \dots, d]) = (n \leq n' \wedge n' + d' \leq n + d)$

$\langle proof \rangle$

lemma *iIN-iFROM-subset-conv*: $([n' \dots, d'] \subseteq [n \dots]) = (n \leq n')$

$\langle proof \rangle$

lemma *iIN-iTILL-subset-conv*: $([n' \dots, d'] \subseteq [\dots n]) = (n' + d' \leq n)$

$\langle proof \rangle$

lemma *iIN-iMOD-subset-conv*:

$$0 < d' \implies ([n' \dots, d'] \subseteq [r, \text{mod } m]) = (r \leq n' \wedge m = \text{Suc } 0)$$

$\langle proof \rangle$

lemma *iIN-iMODb-subset-conv*:

$$0 < d' \implies \\ ([n' \dots, d'] \subseteq [r, \text{mod } m, c]) = \\ (r \leq n' \wedge m = \text{Suc } 0 \wedge n' + d' \leq r + m * c)$$

$\langle proof \rangle$

lemma *iTILL-subset-conv*: $([\dots n'] \subseteq [\dots n]) = (n' \leq n)$

$\langle proof \rangle$

lemma *iTILL-iFROM-subset-conv*: $([\dots n'] \subseteq [n \dots]) = (n = 0)$

$\langle proof \rangle$

lemma *iTILL-iIN-subset-conv*: $([\dots n'] \subseteq [n \dots, d]) = (n = 0 \wedge n' \leq d)$

$\langle proof \rangle$

lemma *iTILL-iMOD-subset-conv*:

$$0 < n' \implies ([\dots n'] \subseteq [r, \text{mod } m]) = (r = 0 \wedge m = \text{Suc } 0)$$

<proof>

lemma *iTILL-iMODb-subset-conv*:

$$0 < n' \implies ([\dots n'] \subseteq [r, \text{mod } m, c]) = (r = 0 \wedge m = \text{Suc } 0 \wedge n' \leq r + m * c)$$

<proof>

lemma *iMOD-iFROM-subset-conv*: $([r', \text{mod } m'] \subseteq [n..]) = (n \leq r')$

<proof>

lemma *iMODb-iFROM-subset-conv*: $([r', \text{mod } m', c'] \subseteq [n..]) = (n \leq r')$

<proof>

lemma *iMODb-iIN-subset-conv*:

$$([r', \text{mod } m', c'] \subseteq [n.., d]) = (n \leq r' \wedge r' + m' * c' \leq n + d)$$

<proof>

lemma *iMODb-iTILL-subset-conv*:

$$([r', \text{mod } m', c'] \subseteq [\dots n]) = (r' + m' * c' \leq n)$$

<proof>

lemma *iMOD-0-subset-conv*: $([r', \text{mod } 0] \subseteq [r, \text{mod } m]) = (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r')$

<proof>

lemma *iMOD-subset-conv*: $0 < m \implies$

$$([r', \text{mod } m'] \subseteq [r, \text{mod } m]) = (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0)$$

<proof>

lemma *iMODb-subset-mod-0-conv*:

$$([r', \text{mod } m', c'] \subseteq [r, \text{mod } 0, c]) = (r'=r \wedge (m'=0 \vee c'=0))$$

<proof>

lemma *iMODb-subset-0-conv*:

$$([r', \text{mod } m', c'] \subseteq [r, \text{mod } m, 0]) = (r'=r \wedge (m'=0 \vee c'=0))$$

<proof>

lemma *iMODb-0-subset-conv*:

$$([r', \text{mod } m', 0] \subseteq [r, \text{mod } m, c]) = (r' \in [r, \text{mod } m, c])$$

<proof>

lemma *iMODb-mod-0-subset-conv*:

$$([r', \text{mod } 0, c'] \subseteq [r, \text{mod } m, c]) = (r' \in [r, \text{mod } m, c])$$

<proof>

lemma *iMODb-subset-conv'*: $\llbracket 0 < c; 0 < c' \rrbracket \implies$

$$([r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]) = (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0 \wedge$$

$$r' + m' * c' \leq r + m * c)$$

$\langle \text{proof} \rangle$

lemma *iMODb-subset-conv*: $\llbracket 0 < m'; 0 < c' \rrbracket \implies$
 $([r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]) =$
 $(r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0 \wedge$
 $r' + m' * c' \leq r + m * c)$
 $\langle \text{proof} \rangle$

lemma *iMODb-iMOD-subset-conv*: $0 < c' \implies$
 $([r', \text{mod } m', c'] \subseteq [r, \text{mod } m]) =$
 $(r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0)$
 $\langle \text{proof} \rangle$

lemmas *iT-subset-conv* =
iFROM-subset-conv
iFROM-iMOD-subset-conv
iTILL-subset-conv
iTILL-iFROM-subset-conv
iTILL-iIN-subset-conv
iTILL-iMOD-subset-conv
iTILL-iMODb-subset-conv
iIN-subset-conv
iIN-iFROM-subset-conv
iIN-iTILL-subset-conv
iIN-iMOD-subset-conv
iIN-iMODb-subset-conv
iMOD-subset-conv
iMOD-iFROM-subset-conv
iMODb-subset-conv'
iMODb-subset-conv
iMODb-iFROM-subset-conv
iMODb-iIN-subset-conv
iMODb-iTILL-subset-conv
iMODb-iMOD-subset-conv

lemma *iFROM-subset*: $n \leq n' \implies [n' \dots] \subseteq [n \dots]$
 $\langle \text{proof} \rangle$

lemma *not-iFROM-iIN-subset*: $\neg [n' \dots] \subseteq [n \dots, d]$
 $\langle \text{proof} \rangle$

lemma *not-iFROM-iTILL-subset*: $\neg [n' \dots] \subseteq [\dots n]$
 $\langle \text{proof} \rangle$

lemma *not-iFROM-iMOD-subset*: $m \neq \text{Suc } 0 \implies \neg [n' \dots] \subseteq [r, \text{mod } m]$
 $\langle \text{proof} \rangle$

lemma *not-iFROM-iMODb-subset*: $\neg [n' \dots] \subseteq [r, \text{mod } m, c]$

$\langle proof \rangle$

lemma *iIN-subset*: $\llbracket n \leq n'; n' + d' \leq n + d \rrbracket \implies [n'..d'] \subseteq [n..d]$
 $\langle proof \rangle$

lemma *iIN-iFROM-subset*: $n \leq n' \implies [n'..d'] \subseteq [n..]$
 $\langle proof \rangle$

lemma *iIN-iTILL-subset*: $n' + d' \leq n \implies [n'..d'] \subseteq [..n]$
 $\langle proof \rangle$

lemma *not-iIN-iMODb-subset*: $\llbracket 0 < d'; m \neq \text{Suc } 0 \rrbracket \implies \neg [n'..d'] \subseteq [r, \text{mod } m, c]$
 $\langle proof \rangle$

lemma *not-iIN-iMOD-subset*: $\llbracket 0 < d'; m \neq \text{Suc } 0 \rrbracket \implies \neg [n'..d'] \subseteq [r, \text{mod } m]$
 $\langle proof \rangle$

lemma *iTILL-subset*: $n' \leq n \implies [..n'] \subseteq [..n]$
 $\langle proof \rangle$

lemma *iTILL-iFROM-subset*: $([..n'] \subseteq [0..])$
 $\langle proof \rangle$

lemma *iTILL-iIN-subset*: $n' \leq d \implies ([..n'] \subseteq [0..d])$
 $\langle proof \rangle$

lemma *not-iTILL-iMOD-subset*:
 $\llbracket 0 < n'; m \neq \text{Suc } 0 \rrbracket \implies \neg [..n'] \subseteq [r, \text{mod } m]$
 $\langle proof \rangle$

lemma *not-iTILL-iMODb-subset*:
 $\llbracket 0 < n'; m \neq \text{Suc } 0 \rrbracket \implies \neg [..n'] \subseteq [r, \text{mod } m, c]$
 $\langle proof \rangle$

lemma *iMOD-iFROM-subset*: $n \leq r' \implies [r', \text{mod } m'] \subseteq [n..]$
 $\langle proof \rangle$

lemma *not-iMOD-iIN-subset*: $0 < m' \implies \neg [r', \text{mod } m'] \subseteq [n..d]$
 $\langle proof \rangle$

lemma *not-iMOD-iTILL-subset*: $0 < m' \implies \neg [r', \text{mod } m'] \subseteq [..n]$
 $\langle proof \rangle$

lemma *iMOD-subset*:
 $\llbracket r \leq r'; r' \text{ mod } m = r \text{ mod } m; m' \text{ mod } m = 0 \rrbracket \implies [r', \text{mod } m'] \subseteq [r, \text{mod } m]$
 $\langle proof \rangle$

lemma *not-iMOD-iMODb-subset*: $0 < m' \implies \neg [r', \text{mod } m'] \subseteq [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

lemma *iMODb-iFROM-subset*: $n \leq r' \implies [r', \text{mod } m', c'] \subseteq [n..]$
 $\langle \text{proof} \rangle$

lemma *iMODb-iTILL-subset*:
 $r' + m' * c' \leq n \implies [r', \text{mod } m', c'] \subseteq [..n]$
 $\langle \text{proof} \rangle$

lemma *iMODb-iIN-subset*:
 $\llbracket n \leq r'; r' + m' * c' \leq n + d \rrbracket \implies [r', \text{mod } m', c'] \subseteq [n..,d]$
 $\langle \text{proof} \rangle$

lemma *iMODb-iMOD-subset*:
 $\llbracket r \leq r'; r' \text{ mod } m = r \text{ mod } m; m' \text{ mod } m = 0 \rrbracket \implies [r', \text{mod } m', c'] \subseteq [r, \text{mod } m]$
 $\langle \text{proof} \rangle$

lemma *iMODb-subset*:
 $\llbracket r \leq r'; r' \text{ mod } m = r \text{ mod } m; m' \text{ mod } m = 0; r' + m' * c' \leq r + m * c \rrbracket \implies$
 $[r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

lemma *iFROM-trans*: $\llbracket y \in [x..]; z \in [y..] \rrbracket \implies z \in [x..]$
 $\langle \text{proof} \rangle$

lemma *iTILL-trans*: $\llbracket y \in [..x]; z \in [..y] \rrbracket \implies z \in [..x]$
 $\langle \text{proof} \rangle$

lemma *iIN-trans*:
 $\llbracket y \in [x..,d]; z \in [y..,d']; d' \leq x + d - y \rrbracket \implies z \in [x..,d]$
 $\langle \text{proof} \rangle$

lemma *iMOD-trans*:
 $\llbracket y \in [x, \text{mod } m]; z \in [y, \text{mod } m] \rrbracket \implies z \in [x, \text{mod } m]$
 $\langle \text{proof} \rangle$

lemma *iMODb-trans*:
 $\llbracket y \in [x, \text{mod } m, c]; z \in [y, \text{mod } m, c']; m * c' \leq x + m * c - y \rrbracket \implies$
 $z \in [x, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

lemma *iMODb-trans'*:
 $\llbracket y \in [x, \text{mod } m, c]; z \in [y, \text{mod } m, c']; c' \leq x \text{ div } m + c - y \text{ div } m \rrbracket \implies$
 $z \in [x, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

1.3.3 Equality of intervals

lemma *iFROM-eq-conv*: $([n..] = [n'..]) = (n = n')$
 $\langle proof \rangle$

lemma *iIN-eq-conv*: $([n..,d] = [n'..,d']) = (n = n' \wedge d = d')$
 $\langle proof \rangle$

lemma *iTILL-eq-conv*: $([...\,n] = [...\,n']) = (n = n')$
 $\langle proof \rangle$

lemma *iMOD-0-eq-conv*: $([r, \text{mod } 0] = [r', \text{mod } m']) = (r = r' \wedge m' = 0)$
 $\langle proof \rangle$

lemma *iMOD-eq-conv*: $0 < m \implies ([r, \text{mod } m] = [r', \text{mod } m']) = (r = r' \wedge m = m')$
 $\langle proof \rangle$

lemma *iMODb-mod-0-eq-conv*:
 $([r, \text{mod } 0, c] = [r', \text{mod } m', c']) = (r = r' \wedge (m' = 0 \vee c' = 0))$
 $\langle proof \rangle$

lemma *iMODb-0-eq-conv*:
 $([r, \text{mod } m, 0] = [r', \text{mod } m', c']) = (r = r' \wedge (m' = 0 \vee c' = 0))$
 $\langle proof \rangle$

lemma *iMODb-eq-conv*: $\llbracket 0 < m; 0 < c \rrbracket \implies$
 $([r, \text{mod } m, c] = [r', \text{mod } m', c']) = (r = r' \wedge m = m' \wedge c = c')$
 $\langle proof \rangle$

lemma *iMOD-iFROM-eq-conv*: $([n..] = [r, \text{mod } m]) = (n = r \wedge m = \text{Suc } 0)$
 $\langle proof \rangle$

lemma *iMODb-iIN-0-eq-conv*:
 $([n..,0] = [r, \text{mod } m, c]) = (n = r \wedge (m = 0 \vee c = 0))$
 $\langle proof \rangle$

lemma *iMODb-iIN-eq-conv*:
 $0 < d \implies ([n..,d] = [r, \text{mod } m, c]) = (n = r \wedge m = \text{Suc } 0 \wedge c = d)$
 $\langle proof \rangle$

1.3.4 Inequality of intervals

lemma *iFROM-iIN-neg*: $[n'..] \neq [n..,d]$
 $\langle proof \rangle$

corollary *iFROM-iTILL-neg*: $[n'..] \neq [...\,n]$
 $\langle proof \rangle$

corollary *iFROM-iMOD-neg*: $m \neq \text{Suc } 0 \implies [n..] \neq [r, \text{mod } m]$

$\langle proof \rangle$

corollary *iFROM-iMODb-neq*: $[n\dots] \neq [r, \text{mod } m, c]$

$\langle proof \rangle$

corollary *iIN-iMOD-neq*: $0 < m \implies [n\dots, d] \neq [r, \text{mod } m]$

$\langle proof \rangle$

corollary *iIN-iMODb-neq2*: $\llbracket m \neq \text{Suc } 0; 0 < d \rrbracket \implies [n\dots, d] \neq [r, \text{mod } m, c]$

$\langle proof \rangle$

lemma *iIN-iMODb-neq*: $\llbracket 2 \leq m; 0 < c \rrbracket \implies [n\dots, d] \neq [r, \text{mod } m, c]$

$\langle proof \rangle$

lemma *iTILL-iIN-neq*: $0 < n \implies [\dots n^\frown] \neq [n\dots, d]$

$\langle proof \rangle$

corollary *iTILL-iMOD-neq*: $0 < m \implies [\dots n] \neq [r, \text{mod } m]$

$\langle proof \rangle$

corollary *iTILL-iMODb-neq*:

$\llbracket m \neq \text{Suc } 0; 0 < n \rrbracket \implies [\dots n] \neq [r, \text{mod } m, c]$

$\langle proof \rangle$

lemma *iMOD-iMODb-neq*: $0 < m \implies [r, \text{mod } m] \neq [r', \text{mod } m', c]$

$\langle proof \rangle$

lemmas *iT-neq* =

iFROM-iTILL-neq iFROM-iIN-neq iFROM-iMOD-neq iFROM-iMODb-neq

iTILL-iIN-neq iTILL-iMOD-neq iTILL-iMODb-neq

iIN-iMOD-neq iIN-iMODb-neq iIN-iMODb-neq2

iMOD-iMODb-neq

1.4 Union and intersection of intervals

lemma *iFROM-union'*: $[n\dots] \cup [n'\dots] = [\min n \ n'\dots]$

$\langle proof \rangle$

corollary *iFROM-union*: $n \leq n' \implies [n\dots] \cup [n'\dots] = [n\dots]$

$\langle proof \rangle$

lemma *iFROM-inter'*: $[n\dots] \cap [n'\dots] = [\max n \ n'\dots]$

$\langle proof \rangle$

corollary *iFROM-inter*: $n' \leq n \implies [n\dots] \cap [n'\dots] = [n\dots]$

$\langle proof \rangle$

lemma *iTILL-union'*: $[\dots n] \cup [\dots n^\frown] = [\dots \max n \ n^\frown]$

$\langle proof \rangle$

corollary *iTILL-union*: $n' \leq n \implies [\dots n] \cup [\dots n'] = [\dots n]$
 $\langle \text{proof} \rangle$

lemma *iTILL-iFROM-union*: $n \leq n' \implies [\dots n'] \cup [n\dots] = UNIV$
 $\langle \text{proof} \rangle$

lemma *iTILL-inter'*: $[\dots n] \cap [\dots n'] = [\dots \min n n']$
 $\langle \text{proof} \rangle$

corollary *iTILL-inter*: $n \leq n' \implies [\dots n] \cap [\dots n'] = [\dots n]$
 $\langle \text{proof} \rangle$

Union and intersection for iIN only when there intersection is not empty and none of them is other's subset, for instance: $\dots n \dots n+d \dots n' \dots n'+d'$

lemma *iIN-union*:
 $\llbracket n \leq n'; n' \leq \text{Suc } (n + d); n + d \leq n' + d' \rrbracket \implies$
 $[n\dots, d] \cup [n'\dots, d'] = [n\dots, n' - n + d']$
 $\langle \text{proof} \rangle$

lemma *iIN-append-union*:
 $[n\dots, d] \cup [n + d\dots, d'] = [n\dots, d + d']$
 $\langle \text{proof} \rangle$

lemma *iIN-append-union-Suc*:
 $[n\dots, d] \cup [\text{Suc } (n + d)\dots, d'] = [n\dots, \text{Suc } (d + d')]$
 $\langle \text{proof} \rangle$

lemma *iIN-append-union-pred*:
 $0 < d \implies [n\dots, d - \text{Suc } 0] \cup [n + d\dots, d'] = [n\dots, d + d']$
 $\langle \text{proof} \rangle$

lemma *iIN-iFROM-union*:
 $n' \leq \text{Suc } (n + d) \implies [n\dots, d] \cup [n'\dots] = [\min n n'\dots]$
 $\langle \text{proof} \rangle$

lemma *iIN-iFROM-append-union*:
 $[n\dots, d] \cup [n + d\dots] = [n\dots]$
 $\langle \text{proof} \rangle$

lemma *iIN-iFROM-append-union-Suc*:
 $[n\dots, d] \cup [\text{Suc } (n + d)\dots] = [n\dots]$
 $\langle \text{proof} \rangle$

lemma *iIN-iFROM-append-union-pred*:
 $0 < d \implies [n\dots, d - \text{Suc } 0] \cup [n + d\dots] = [n\dots]$
 $\langle \text{proof} \rangle$

lemma *iIN-inter*:

$$\begin{aligned} & \llbracket n \leq n'; n' \leq n + d; n + d \leq n' + d' \rrbracket \implies \\ & [n \dots, d] \cap [n' \dots, d'] = [n' \dots, n + d - n'] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMOD-union*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m \rrbracket \implies \\ & [r, \bmod m] \cup [r', \bmod m] = [r, \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMOD-union'*:

$$\begin{aligned} & r \bmod m = r' \bmod m \implies \\ & [r, \bmod m] \cup [r', \bmod m] = [\min r \ r', \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMOD-inter*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m \rrbracket \implies \\ & [r, \bmod m] \cap [r', \bmod m] = [r', \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMOD-inter'*:

$$\begin{aligned} & r \bmod m = r' \bmod m \implies \\ & [r, \bmod m] \cap [r', \bmod m] = [\max r \ r', \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-union*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m; r' \leq r + m * c; r + m * c \leq r' + m * c' \rrbracket \implies \\ & [r, \bmod m, c] \cup [r', \bmod m, c'] = [r, \bmod m, r' \text{ div } m - r \text{ div } m + c'] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-append-union*:

$$\begin{aligned} & [r, \bmod m, c] \cup [r + m * c, \bmod m, c'] = [r, \bmod m, c + c'] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-iMOD-append-union'*:

$$\begin{aligned} & \llbracket r \bmod m = r' \bmod m; r' \leq r + m * \text{Suc } c \rrbracket \implies \\ & [r, \bmod m, c] \cup [r', \bmod m] = [\min r \ r', \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-iMOD-append-union*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m; r' \leq r + m * \text{Suc } c \rrbracket \implies \\ & [r, \bmod m, c] \cup [r', \bmod m] = [r, \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-append-union-Suc*:

$$\begin{aligned} & [r, \bmod m, c] \cup [r + m * \text{Suc } c, \bmod m, c'] = [r, \bmod m, \text{Suc } (c + c')] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-append-union-pred*:

$$0 < c \implies [r, \bmod m, c - \text{Suc } 0] \cup [r + m * c, \bmod m, c'] = [r, \bmod m, c + c']$$

$c']$
 $\langle proof \rangle$

lemma *iMODb-inter*:

$\llbracket r \leq r'; r \bmod m = r' \bmod m; r' \leq r + m * c; r + m * c \leq r' + m * c' \rrbracket \implies$
 $[r, \bmod m, c] \cap [r', \bmod m, c'] = [r', \bmod m, c - (r' - r) \text{ div } m]$
 $\langle proof \rangle$

lemmas *iT-union'* =

iFROM-union'
iTILL-union'
iMOD-union'
iMODb-iMOD-append-union'

lemmas *iT-union* =

iFROM-union
iTILL-union
iTILL-iFROM-union
iIN-union
iIN-iFROM-union
iMOD-union
iMODb-union

lemmas *iT-union-append* =

iIN-append-union
iIN-append-union-Suc
iIN-append-union-pred
iIN-iFROM-append-union
iIN-iFROM-append-union-Suc
iIN-iFROM-append-union-pred
iMODb-append-union
iMODb-iMOD-append-union
iMODb-append-union-Suc
iMODb-append-union-pred

lemmas *iT-inter'* =

iFROM-inter'
iTILL-inter'
iMOD-inter'

lemmas *iT-inter* =

iFROM-inter
iTILL-inter
iIN-inter
iMOD-inter
iMODb-inter

lemma *mod-partition-Union*:

$0 < m \implies (\bigcup k. A \cap [k * m \dots m - \text{Suc } 0]) = A$

$\langle proof \rangle$

lemma *finite-mod-partition-Union*:

$\llbracket 0 < m; \text{finite } A \rrbracket \implies$
 $(\bigcup_{k \leq \text{Max } A \text{ div } m} A \cap [k * m \dots, m - \text{Suc } 0]) = A$
 $\langle proof \rangle$

lemma *mod-partition-is-disjoint*:

$\llbracket 0 < (m :: \text{nat}); k \neq k' \rrbracket \implies$
 $(A \cap [k * m \dots, m - \text{Suc } 0]) \cap (A \cap [k' * m \dots, m - \text{Suc } 0]) = \{\}$
 $\langle proof \rangle$

1.5 Cutting intervals

lemma *iTILL-cut-le*: $[\dots n] \downarrow \leq t = (\text{if } t \leq n \text{ then } [\dots t] \text{ else } [\dots n])$
 $\langle proof \rangle$

corollary *iTILL-cut-le1*: $t \in [\dots n] \implies [\dots n] \downarrow \leq t = [\dots t]$
 $\langle proof \rangle$

corollary *iTILL-cut-le2*: $t \notin [\dots n] \implies [\dots n] \downarrow \leq t = [\dots n]$
 $\langle proof \rangle$

lemma *iFROM-cut-le*:

$[n \dots] \downarrow \leq t =$
 $(\text{if } t < n \text{ then } \{\} \text{ else } [n \dots, t - n])$
 $\langle proof \rangle$

corollary *iFROM-cut-le1*: $t \in [n \dots] \implies [n \dots] \downarrow \leq t = [n \dots, t - n]$
 $\langle proof \rangle$

lemma *iIN-cut-le*:

$[n \dots, d] \downarrow \leq t =$
 $(\text{if } t < n \text{ then } \{\} \text{ else}$
 $\text{if } t \leq n + d \text{ then } [n \dots, t - n]$
 $\text{else } [n \dots, d])$
 $\langle proof \rangle$

corollary *iIN-cut-le1*:

$t \in [n \dots, d] \implies [n \dots, d] \downarrow \leq t = [n \dots, t - n]$
 $\langle proof \rangle$

lemma *iMOD-cut-le*:

$[r, \text{mod } m] \downarrow \leq t =$
 $(\text{if } t < r \text{ then } \{\}$
 $\text{else } [r, \text{mod } m, (t - r) \text{ div } m])$
 $\langle proof \rangle$

lemma *iMOD-cut-le1*:

$t \in [r, \text{mod } m] \implies$
 $[r, \text{mod } m] \downarrow \leq t = [r, \text{mod } m, (t - r) \text{ div } m]$
 $\langle \text{proof} \rangle$

lemma *iMODb-cut-le*:

$[r, \text{mod } m, c] \downarrow \leq t =$ (
 if $t < r$ then $\{\}$
 else if $t < r + m * c$ then $[r, \text{mod } m, (t - r) \text{ div } m]$
 else $[r, \text{mod } m, c]$)
 $\langle \text{proof} \rangle$

lemma *iMODb-cut-le1*:

$t \in [r, \text{mod } m, c] \implies$
 $[r, \text{mod } m, c] \downarrow \leq t = [r, \text{mod } m, (t - r) \text{ div } m]$
 $\langle \text{proof} \rangle$

lemma *iTILL-cut-less*:

$[\dots n] \downarrow < t =$ (
 if $n < t$ then $[\dots n]$ else
 if $t = 0$ then $\{\}$
 else $[\dots t - \text{Suc } 0]$)
 $\langle \text{proof} \rangle$

lemma *iTILL-cut-less1*:

$\llbracket t \in [\dots n]; 0 < t \rrbracket \implies [\dots n] \downarrow < t = [\dots t - \text{Suc } 0]$
 $\langle \text{proof} \rangle$

lemma *iFROM-cut-less*:

$[n \dots] \downarrow < t =$ (
 if $t \leq n$ then $\{\}$
 else $[n \dots, t - \text{Suc } n]$)
 $\langle \text{proof} \rangle$

lemma *iFROM-cut-less1*:

$n < t \implies [n \dots] \downarrow < t = [n \dots, t - \text{Suc } n]$
 $\langle \text{proof} \rangle$

lemma *iIN-cut-less*:

$[n \dots, d] \downarrow < t =$ (
 if $t \leq n$ then $\{\}$ else
 if $t \leq n + d$ then $[n \dots, t - \text{Suc } n]$
 else $[n \dots, d]$)
 $\langle \text{proof} \rangle$

lemma *iIN-cut-less1*:

$\llbracket t \in [n \dots, d]; n < t \rrbracket \implies [n \dots, d] \downarrow < t = [n \dots, t - \text{Suc } n]$
 $\langle \text{proof} \rangle$

lemma *iMOD-cut-less*:

$[r, \text{mod } m] \downarrow < t = ($
 $\quad \text{if } t \leq r \text{ then } \{\}$
 $\quad \text{else } [r, \text{mod } m, (t - \text{Suc } r) \text{ div } m])$
 $\langle \text{proof} \rangle$

lemma *iMOD-cut-less1*:

$\llbracket t \in [r, \text{mod } m]; r < t \rrbracket \implies$
 $[r, \text{mod } m] \downarrow < t = [r, \text{mod } m, (t - r) \text{ div } m - \text{Suc } 0]$
 $\langle \text{proof} \rangle$

lemma *iMODb-cut-less*:

$[r, \text{mod } m, c] \downarrow < t = ($
 $\quad \text{if } t \leq r \text{ then } \{\} \text{ else}$
 $\quad \text{if } r + m * c < t \text{ then } [r, \text{mod } m, c]$
 $\quad \text{else } [r, \text{mod } m, (t - \text{Suc } r) \text{ div } m])$
 $\langle \text{proof} \rangle$

lemma *iMODb-cut-less1*: $\llbracket t \in [r, \text{mod } m, c]; r < t \rrbracket \implies$

$[r, \text{mod } m, c] \downarrow < t = [r, \text{mod } m, (t - r) \text{ div } m - \text{Suc } 0]$
 $\langle \text{proof} \rangle$

lemmas *iT-cut-le* =

iTILL-cut-le
iFROM-cut-le
iIN-cut-le
iMOD-cut-le
iMODb-cut-le

lemmas *iT-cut-le1* =

iTILL-cut-le1
iFROM-cut-le1
iIN-cut-le1
iMOD-cut-le1
iMODb-cut-le1

lemmas *iT-cut-less* =

iTILL-cut-less
iFROM-cut-less
iIN-cut-less
iMOD-cut-less
iMODb-cut-less

lemmas *iT-cut-less1* =
iTILL-cut-less1
iFROM-cut-less1
iIN-cut-less1
iMOD-cut-less1
iMODb-cut-less1

lemmas *iT-cut-le-less* =
iTILL-cut-le
iTILL-cut-less
iFROM-cut-le
iFROM-cut-less
iIN-cut-le
iIN-cut-less
iMOD-cut-le
iMOD-cut-less
iMODb-cut-le
iMODb-cut-less

lemmas *iT-cut-le-less1* =
iTILL-cut-le1
iTILL-cut-less1
iFROM-cut-le1
iFROM-cut-less1
iIN-cut-le1
iIN-cut-less1
iMOD-cut-le1
iMOD-cut-less1
iMODb-cut-le1
iMODb-cut-less1

lemma *iTILL-cut-ge*:
 $[\dots n] \downarrow \geq t = (\text{if } n < t \text{ then } \{\} \text{ else } [t \dots, n-t])$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-ge1*: $t \in [\dots n] \implies [\dots n] \downarrow \geq t = [t \dots, n-t]$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-ge2*: $t \notin [\dots n] \implies [\dots n] \downarrow \geq t = \{\}$
 $\langle \text{proof} \rangle$

lemma *iTILL-cut-greater*:
 $[\dots n] \downarrow > t = (\text{if } n \leq t \text{ then } \{\} \text{ else } [\text{Suc } t \dots, n - \text{Suc } t])$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-greater1*:
 $t \in [\dots n] \implies t < n \implies [\dots n] \downarrow > t = [\text{Suc } t \dots, n - \text{Suc } t]$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-greater2*: $t \notin [\dots n] \implies [\dots n] \downarrow > t = \{\}$
 $\langle \text{proof} \rangle$

lemma *iFROM-cut-ge*:

$[n\dots] \downarrow \geq t = (\text{if } n \leq t \text{ then } [t\dots] \text{ else } [n\dots])$

$\langle \text{proof} \rangle$

corollary *iFROM-cut-ge1*: $t \in [n\dots] \implies [n\dots] \downarrow \geq t = [t\dots]$

$\langle \text{proof} \rangle$

lemma *iFROM-cut-greater*:

$[n\dots] \downarrow > t = (\text{if } n \leq t \text{ then } [\text{Suc } t\dots] \text{ else } [n\dots])$

$\langle \text{proof} \rangle$

corollary *iFROM-cut-greater1*:

$t \in [n\dots] \implies [n\dots] \downarrow > t = [\text{Suc } t\dots]$

$\langle \text{proof} \rangle$

lemma *iIN-cut-ge*:

$[n\dots, d] \downarrow \geq t = (\text{if } t < n \text{ then } [n\dots, d] \text{ else } \text{if } t \leq n+d \text{ then } [t\dots, n+d-t] \text{ else } \{\})$

$\langle \text{proof} \rangle$

corollary *iIN-cut-ge1*: $t \in [n\dots, d] \implies$

$[n\dots, d] \downarrow \geq t = [t\dots, n+d-t]$

$\langle \text{proof} \rangle$

corollary *iIN-cut-ge2*: $t \notin [n\dots, d] \implies$

$[n\dots, d] \downarrow \geq t = (\text{if } t < n \text{ then } [n\dots, d] \text{ else } \{\})$

$\langle \text{proof} \rangle$

lemma *iIN-cut-greater*:

$[n\dots, d] \downarrow > t = (\text{if } t < n \text{ then } [n\dots, d] \text{ else } \text{if } t < n+d \text{ then } [\text{Suc } t\dots, n+d - \text{Suc } t] \text{ else } \{\})$

$\langle \text{proof} \rangle$

corollary *iIN-cut-greater1*:

$\llbracket t \in [n\dots, d]; t < n+d \rrbracket \implies$

$[n\dots, d] \downarrow > t = [\text{Suc } t\dots, n+d - \text{Suc } t]$

$\langle \text{proof} \rangle$

lemma *mod-cut-greater-aux-t-less*:

$$\llbracket 0 < (m::nat); r \leq t \rrbracket \implies$$

$$t < t + m - (t - r) \bmod m$$

$$\langle proof \rangle$$

lemma *mod-cut-greater-aux-le-x*:

$$\llbracket (r::nat) \leq t; t < x; x \bmod m = r \bmod m \rrbracket \implies$$

$$t + m - (t - r) \bmod m \leq x$$

$$\langle proof \rangle$$

lemma *iMOD-cut-greater*:

$$[r, \bmod m] \downarrow > t = ($$

$$\text{if } t < r \text{ then } [r, \bmod m] \text{ else}$$

$$\text{if } m = 0 \text{ then } \{\}$$

$$\text{else } [t + m - (t - r) \bmod m, \bmod m])$$

$$\langle proof \rangle$$

lemma *iMOD-cut-greater1*:

$$t \in [r, \bmod m] \implies$$

$$[r, \bmod m] \downarrow > t = ($$

$$\text{if } m = 0 \text{ then } \{\}$$

$$\text{else } [t + m, \bmod m])$$

$$\langle proof \rangle$$

lemma *iMODb-cut-greater-aux*:

$$\llbracket 0 < m; t < r + m * c; r \leq t \rrbracket \implies$$

$$(r + m * c - (t + m - (t - r) \bmod m)) \operatorname{div} m =$$

$$c - \operatorname{Suc} ((t - r) \operatorname{div} m)$$

$$\langle proof \rangle$$

lemma *iMODb-cut-greater*:

$$[r, \bmod m, c] \downarrow > t = ($$

$$\text{if } t < r \text{ then } [r, \bmod m, c] \text{ else}$$

$$\text{if } r + m * c \leq t \text{ then } \{\}$$

$$\text{else } [t + m - (t - r) \bmod m, \bmod m, c - \operatorname{Suc} ((t - r) \operatorname{div} m)])$$

$$\langle proof \rangle$$

lemma *iMODb-cut-greater1*:

$$t \in [r, \bmod m, c] \implies$$

$$[r, \bmod m, c] \downarrow > t = ($$

$$\text{if } r + m * c \leq t \text{ then } \{\}$$

$$\text{else } [t + m, \bmod m, c - \operatorname{Suc} ((t - r) \operatorname{div} m)])$$

$$\langle proof \rangle$$

lemma *iMOD-cut-ge*:

$$[r, \bmod m] \downarrow \geq t = ($$

$$\text{if } t \leq r \text{ then } [r, \bmod m] \text{ else}$$

$$\text{if } m = 0 \text{ then } \{\}$$

$\text{else } [t + m - \text{Suc } ((t - \text{Suc } r) \bmod m), \bmod m]$
 $\langle \text{proof} \rangle$

lemma *iMOD-cut-ge1*:
 $t \in [r, \bmod m] \implies$
 $[r, \bmod m] \downarrow \geq t = [t, \bmod m]$
 $\langle \text{proof} \rangle$

lemma *iMODb-cut-ge*:
 $[r, \bmod m, c] \downarrow \geq t =$
 $\text{if } t \leq r \text{ then } [r, \bmod m, c] \text{ else}$
 $\text{if } r + m * c < t \text{ then } \{\}$
 $\text{else } [t + m - \text{Suc } ((t - \text{Suc } r) \bmod m), \bmod m, c - (t + m - \text{Suc } r) \text{ div } m]$
 $\langle \text{proof} \rangle$

lemma *iMODb-cut-ge1*:
 $t \in [r, \bmod m, c] \implies$
 $[r, \bmod m, c] \downarrow \geq t =$
 $\text{if } r + m * c < t \text{ then } \{\}$
 $\text{else } [t, \bmod m, c - (t - r) \text{ div } m]$
 $\langle \text{proof} \rangle$

lemma *iMOD-0-cut-greater*:
 $t \in [r, \bmod 0] \implies [r, \bmod 0] \downarrow > t = \{\}$
 $\langle \text{proof} \rangle$

lemma *iMODb-0-cut-greater*: $t \in [r, \bmod 0, c] \implies$
 $[r, \bmod 0, c] \downarrow > t = \{\}$
 $\langle \text{proof} \rangle$

lemmas *iT-cut-ge* =
iTILL-cut-ge
iFROM-cut-ge
iIN-cut-ge
iMOD-cut-ge
iMODb-cut-ge

lemmas *iT-cut-ge1* =
iTILL-cut-ge1
iFROM-cut-ge1
iIN-cut-ge1
iMOD-cut-ge1
iMODb-cut-ge1

lemmas *iT-cut-greater* =
iTILL-cut-greater
iFROM-cut-greater
iIN-cut-greater

iMOD-cut-greater
iMODb-cut-greater

lemmas *iT-cut-greater1* =
iTILL-cut-greater1
iFROM-cut-greater1
iIN-cut-greater1
iMOD-cut-greater1
iMODb-cut-greater1

lemmas *iT-cut-ge-greater* =
iTILL-cut-ge
iTILL-cut-greater
iFROM-cut-ge
iFROM-cut-greater
iIN-cut-ge
iIN-cut-greater
iMOD-cut-ge
iMOD-cut-greater
iMODb-cut-ge
iMODb-cut-greater

lemmas *iT-cut-ge-greater1* =
iTILL-cut-ge1
iTILL-cut-greater1
iFROM-cut-ge1
iFROM-cut-greater1
iIN-cut-ge1
iIN-cut-greater1
iMOD-cut-ge1
iMOD-cut-greater1
iMODb-cut-ge1
iMODb-cut-greater1

1.6 Cardinality of intervals

lemma *iFROM-card*: $\text{card } [n..] = 0$
 $\langle \text{proof} \rangle$

lemma *iTILL-card*: $\text{card } [...n] = \text{Suc } n$
 $\langle \text{proof} \rangle$

lemma *iIN-card*: $\text{card } [n..., d] = \text{Suc } d$
 $\langle \text{proof} \rangle$

lemma *iMOD-0-card*: $\text{card } [r, \text{mod } 0] = \text{Suc } 0$
 $\langle \text{proof} \rangle$

lemma *iMOD-card*: $0 < m \implies \text{card } [r, \text{mod } m] = 0$
 $\langle \text{proof} \rangle$

lemma *iMOD-card-if*: $\text{card } [r, \text{mod } m] = (\text{if } m = 0 \text{ then } \text{Suc } 0 \text{ else } 0)$
 $\langle \text{proof} \rangle$

lemma *iMODb-mod-0-card*: $\text{card } [r, \text{mod } 0, c] = \text{Suc } 0$
 $\langle \text{proof} \rangle$

lemma *iMODb-card*: $0 < m \implies \text{card } [r, \text{mod } m, c] = \text{Suc } c$
 $\langle \text{proof} \rangle$

lemma *iMODb-card-if*:
 $\text{card } [r, \text{mod } m, c] = (\text{if } m = 0 \text{ then } \text{Suc } 0 \text{ else } \text{Suc } c)$
 $\langle \text{proof} \rangle$

lemmas *iT-card* =
iFROM-card
iTILL-card
iIN-card
iMOD-card-if
iMODb-card-if

Cardinality with *icard*

lemma *iFROM-icard*: $\text{icard } [n..] = \infty$
 $\langle \text{proof} \rangle$

lemma *iTILL-icard*: $\text{icard } [..n] = \text{enat } (\text{Suc } n)$
 $\langle \text{proof} \rangle$

lemma *iIN-icard*: $\text{icard } [n.., d] = \text{enat } (\text{Suc } d)$
 $\langle \text{proof} \rangle$

lemma *iMOD-0-icard*: $\text{icard } [r, \text{mod } 0] = \text{eSuc } 0$
 $\langle \text{proof} \rangle$

lemma *iMOD-icard*: $0 < m \implies \text{icard } [r, \text{mod } m] = \infty$
 $\langle \text{proof} \rangle$

lemma *iMOD-icard-if*: $\text{icard } [r, \text{mod } m] = (\text{if } m = 0 \text{ then } \text{eSuc } 0 \text{ else } \infty)$
 $\langle \text{proof} \rangle$

lemma *iMODb-mod-0-icard*: $\text{icard } [r, \text{mod } 0, c] = \text{eSuc } 0$
 $\langle \text{proof} \rangle$

lemma *iMODb-icard*: $0 < m \implies \text{icard } [r, \text{mod } m, c] = \text{enat } (\text{Suc } c)$
 $\langle \text{proof} \rangle$

lemma *iMODb-icard-if*: $\text{icard } [r, \text{mod } m, c] = \text{enat } (\text{if } m = 0 \text{ then } \text{Suc } 0 \text{ else } \text{Suc } c)$
 $\langle \text{proof} \rangle$

lemmas *iT-icard* =
iFROM-icard
iTILL-icard
iIN-icard
iMOD-icard-if
iMODb-icard-if

1.7 Functions *inext* and *iprev* with intervals

lemma

iFROM-inext: $t \in [n..] \implies \text{inext } t [n..] = \text{Suc } t$ **and**
iTILL-inext: $t < n \implies \text{inext } t [..n] = \text{Suc } t$ **and**
iIN-inext: $\llbracket n \leq t; t < n + d \rrbracket \implies \text{inext } t [n..,d] = \text{Suc } t$
 $\langle \text{proof} \rangle$

lemma

iFROM-iprev': $t \in [n..] \implies \text{iprev } (\text{Suc } t) [n..] = t$ **and**
iFROM-iprev: $n < t \implies \text{iprev } t [n..] = t - \text{Suc } 0$ **and**
iTILL-iprev: $t \in [..n] \implies \text{iprev } t [..n] = t - \text{Suc } 0$ **and**
iIN-iprev: $\llbracket n < t; t \leq n + d \rrbracket \implies \text{iprev } t [n..,d] = t - \text{Suc } 0$ **and**
iIN-iprev': $\llbracket n \leq t; t < n + d \rrbracket \implies \text{iprev } (\text{Suc } t) [n..,d] = t$
 $\langle \text{proof} \rangle$

lemma *iMOD-inext*: $t \in [r, \text{mod } m] \implies \text{inext } t [r, \text{mod } m] = t + m$
 $\langle \text{proof} \rangle$

lemma *iMOD-iprev*: $\llbracket t \in [r, \text{mod } m]; r < t \rrbracket \implies \text{iprev } t [r, \text{mod } m] = t - m$
 $\langle \text{proof} \rangle$

lemma *iMOD-iprev'*: $t \in [r, \text{mod } m] \implies \text{iprev } (t + m) [r, \text{mod } m] = t$
 $\langle \text{proof} \rangle$

lemma *iMODb-inext*:

$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies$
 $\text{inext } t [r, \text{mod } m, c] = t + m$
 $\langle \text{proof} \rangle$

lemma *iMODb-iprev*:

$\llbracket t \in [r, \text{mod } m, c]; r < t \rrbracket \implies$
 $\text{iprev } t [r, \text{mod } m, c] = t - m$
 $\langle \text{proof} \rangle$

lemma *iMODb-iprev'*:

$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies$
 $\text{iprev } (t + m) [r, \text{mod } m, c] = t$
 $\langle \text{proof} \rangle$

lemmas *iT-inext* =
iFROM-inext

iTILL-inext
iIN-inext
iMOD-inext
iMODb-inext

lemmas *iT-iprev* =

iFROM-iprev'
iFROM-iprev
iTILL-iprev
iIN-iprev
iIN-iprev'
iMOD-iprev
iMOD-iprev'
iMODb-iprev
iMODb-iprev'

lemma *iFROM-inext-if*:

inext t [n..] = (if t ∈ [n..] then Suc t else t)
 ⟨proof⟩

lemma *iTILL-inext-if*:

inext t [...n] = (if t < n then Suc t else t)
 ⟨proof⟩

lemma *iIN-inext-if*:

inext t [n...d] = (if n ≤ t ∧ t < n + d then Suc t else t)
 ⟨proof⟩

lemma *iMOD-inext-if*:

inext t [r, mod m] = (if t ∈ [r, mod m] then t + m else t)
 ⟨proof⟩

lemma *iMODb-inext-if*:

inext t [r, mod m, c] =
*(if t ∈ [r, mod m, c] ∧ t < r + m * c then t + m else t)*
 ⟨proof⟩

lemmas *iT-inext-if* =

iFROM-inext-if
iTILL-inext-if
iIN-inext-if
iMOD-inext-if
iMODb-inext-if

lemma *iFROM-iprev-if*:

iprev t [n..] = (if n < t then t - Suc 0 else t)
 ⟨proof⟩

lemma *iTILL-iprev-if*:

iprev t [...n] = (if t ∈ [...n] then t - Suc 0 else t)
 ⟨proof⟩

lemma *iIN-iprev-if*:

$iprev\ t\ [n\dots,d] = (if\ n < t \wedge t \leq n + d\ then\ t - Suc\ 0\ else\ t)$
 $\langle proof \rangle$

lemma *iMOD-iprev-if*:

$iprev\ t\ [r,\ mod\ m] =$
 $(if\ t \in [r,\ mod\ m] \wedge r < t\ then\ t - m\ else\ t)$
 $\langle proof \rangle$

lemma *iMODb-iprev-if*:

$iprev\ t\ [r,\ mod\ m,\ c] =$
 $(if\ t \in [r,\ mod\ m,\ c] \wedge r < t\ then\ t - m\ else\ t)$
 $\langle proof \rangle$

lemmas *iT-iprev-if* =

iFROM-iprev-if
iTILL-iprev-if
iIN-iprev-if
iMOD-iprev-if
iMODb-iprev-if

The difference between an element and the next/previous element is constant if the element is different from Min/Max of the interval

lemma *iFROM-inext-diff-const*:

$t \in [n\dots] \implies inext\ t\ [n\dots] - t = Suc\ 0$
 $\langle proof \rangle$

lemma *iFROM-iprev-diff-const*:

$n < t \implies t - iprev\ t\ [n\dots] = Suc\ 0$
 $\langle proof \rangle$

lemma *iFROM-iprev-diff-const'*:

$t \in [n\dots] \implies Suc\ t - iprev\ (Suc\ t)\ [n\dots] = Suc\ 0$
 $\langle proof \rangle$

lemma *iTILL-inext-diff-const*:

$t < n \implies inext\ t\ [\dots n] - t = Suc\ 0$
 $\langle proof \rangle$

lemma *iTILL-iprev-diff-const*:

$\llbracket t \in [\dots n];\ 0 < t \rrbracket \implies t - iprev\ t\ [\dots n] = Suc\ 0$
 $\langle proof \rangle$

lemma *iIN-inext-diff-const*:

$\llbracket n \leq t;\ t < n + d \rrbracket \implies inext\ t\ [n\dots,d] - t = Suc\ 0$
 $\langle proof \rangle$

lemma *iIN-iprev-diff-const*:

$\llbracket n < t;\ t \leq n + d \rrbracket \implies t - iprev\ t\ [n\dots,d] = Suc\ 0$
 $\langle proof \rangle$

lemma *iIN-iprev-diff-const'*:

$\llbracket n \leq t;\ t < n + d \rrbracket \implies Suc\ t - iprev\ (Suc\ t)\ [n\dots,d] = Suc\ 0$

$\langle proof \rangle$

lemma *iMOD-inext-diff-const*:

$$t \in [r, \text{mod } m] \implies \text{inext } t [r, \text{mod } m] - t = m$$

$\langle proof \rangle$

lemma *iMOD-iprev-diff-const'*:

$$t \in [r, \text{mod } m] \implies (t + m) - \text{iprev } (t + m) [r, \text{mod } m] = m$$

$\langle proof \rangle$

lemma *iMOD-iprev-diff-const*:

$$\llbracket t \in [r, \text{mod } m]; r < t \rrbracket \implies t - \text{iprev } t [r, \text{mod } m] = m$$

$\langle proof \rangle$

lemma *iMODb-inext-diff-const*:

$$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies \text{inext } t [r, \text{mod } m, c] - t = m$$

$\langle proof \rangle$

lemma *iMODb-iprev-diff-const'*:

$$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies (t + m) - \text{iprev } (t + m) [r, \text{mod } m, c] = m$$

$\langle proof \rangle$

lemma *iMODb-iprev-diff-const*:

$$\llbracket t \in [r, \text{mod } m, c]; r < t \rrbracket \implies t - \text{iprev } t [r, \text{mod } m, c] = m$$

$\langle proof \rangle$

lemmas *iT-inext-diff-const* =

iFROM-inext-diff-const

iTILL-inext-diff-const

iIN-inext-diff-const

iMOD-inext-diff-const

iMODb-inext-diff-const

lemmas *iT-iprev-diff-const* =

iFROM-iprev-diff-const

iFROM-iprev-diff-const'

iTILL-iprev-diff-const

iIN-iprev-diff-const

iIN-iprev-diff-const'

iMOD-iprev-diff-const'

iMOD-iprev-diff-const

iMODb-iprev-diff-const'

iMODb-iprev-diff-const

1.7.1 Mirroring of intervals

lemma

iIN-mirror-elem: *mirror-elem* $x [n..d] = n + d - x$ **and**

iTILL-mirror-elem: *mirror-elem* $x [..n] = n - x$ **and**

iMODb-mirror-elim: $\text{mirror-elim } x [r, \text{mod } m, c] = r + r + m * c - x$
 $\langle \text{proof} \rangle$

lemma *iMODb-imirror-bounds*:

$r' + m' * c' \leq l + r \implies$
 $\text{imirror-bounds } [r', \text{mod } m', c'] \text{ } l \text{ } r = [l + r - r' - m' * c', \text{mod } m', c']$
 $\langle \text{proof} \rangle$

lemma *iIN-imirror-bounds*:

$n + d \leq l + r \implies \text{imirror-bounds } [n \dots, d] \text{ } l \text{ } r = [l + r - n - d \dots, d]$
 $\langle \text{proof} \rangle$

lemma *iTILL-imirror-bounds*:

$n \leq l + r \implies \text{imirror-bounds } [\dots n] \text{ } l \text{ } r = [l + r - n \dots, n]$
 $\langle \text{proof} \rangle$

lemmas *iT-imirror-bounds* =

iTILL-imirror-bounds
iIN-imirror-bounds
iMODb-imirror-bounds

lemma *iMODb-imirror-ident*: $\text{imirror } [r, \text{mod } m, c] = [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

lemma *iIN-imirror-ident*: $\text{imirror } [n \dots, d] = [n \dots, d]$
 $\langle \text{proof} \rangle$

lemma *iTILL-imirror-ident*: $\text{imirror } [\dots n] = [\dots n]$
 $\langle \text{proof} \rangle$

lemmas *iT-imirror-ident* =

iTILL-imirror-ident
iIN-imirror-ident
iMODb-imirror-ident

1.7.2 Functions *inext-nth* and *iprev-nth* on intervals

lemma *iFROM-inext-nth* : $[n \dots] \rightarrow a = n + a$
 $\langle \text{proof} \rangle$

lemma *iIN-inext-nth* : $a \leq d \implies [n \dots, d] \rightarrow a = n + a$
 $\langle \text{proof} \rangle$

lemma *iIN-iprev-nth*: $a \leq d \implies [n \dots, d] \leftarrow a = n + d - a$
 $\langle \text{proof} \rangle$

lemma *iIN-inext-nth-if* :

$[n \dots, d] \rightarrow a = (\text{if } a \leq d \text{ then } n + a \text{ else } n + d)$
 $\langle \text{proof} \rangle$

lemma *iIN-iprev-nth-if*:

$[n..d] \leftarrow a = (\text{if } a \leq d \text{ then } n + d - a \text{ else } n)$
 $\langle \text{proof} \rangle$

lemma *iTILL-inext-nth* : $a \leq n \implies [\dots n] \rightarrow a = a$
 $\langle \text{proof} \rangle$

lemma *iTILL-inext-nth-if* :
 $[\dots n] \rightarrow a = (\text{if } a \leq n \text{ then } a \text{ else } n)$
 $\langle \text{proof} \rangle$

lemma *iTILL-iprev-nth*: $a \leq n \implies [\dots n] \leftarrow a = n - a$
 $\langle \text{proof} \rangle$

lemma *iTILL-iprev-nth-if*:
 $[\dots n] \leftarrow a = (\text{if } a \leq n \text{ then } n - a \text{ else } 0)$
 $\langle \text{proof} \rangle$

lemma *iMOD-inext-nth*: $[r, \text{mod } m] \rightarrow a = r + m * a$
 $\langle \text{proof} \rangle$

lemma *iMODb-inext-nth*: $a \leq c \implies [r, \text{mod } m, c] \rightarrow a = r + m * a$
 $\langle \text{proof} \rangle$

lemma *iMODb-inext-nth-if*:
 $[r, \text{mod } m, c] \rightarrow a = (\text{if } a \leq c \text{ then } r + m * a \text{ else } r + m * c)$
 $\langle \text{proof} \rangle$

lemma *iMODb-iprev-nth*:
 $a \leq c \implies [r, \text{mod } m, c] \leftarrow a = r + m * (c - a)$
 $\langle \text{proof} \rangle$

lemma *iMODb-iprev-nth-if*:
 $[r, \text{mod } m, c] \leftarrow a = (\text{if } a \leq c \text{ then } r + m * (c - a) \text{ else } r)$
 $\langle \text{proof} \rangle$

lemma *iIN-iFROM-inext-nth*:
 $a \leq d \implies [n..d] \rightarrow a = [n..] \rightarrow a$
 $\langle \text{proof} \rangle$

lemma *iIN-iFROM-inext*:
 $a < n + d \implies \text{inext } a [n..d] = \text{inext } a [n..]$
 $\langle \text{proof} \rangle$

lemma *iMOD-iMODb-inext-nth*:
 $a \leq c \implies [r, \text{mod } m, c] \rightarrow a = [r, \text{mod } m] \rightarrow a$
 $\langle \text{proof} \rangle$

lemma *iMOD-iMODb-inext*:

$$a < r + m * c \implies \text{inext } a [r, \text{mod } m, c] = \text{inext } a [r, \text{mod } m]$$

<proof>

lemma *iMOD-inext-nth-Suc-diff*:

$$([r, \text{mod } m] \rightarrow (\text{Suc } n)) - ([r, \text{mod } m] \rightarrow n) = m$$

<proof>

lemma *iMOD-inext-nth-diff*:

$$([r, \text{mod } m] \rightarrow a) - ([r, \text{mod } m] \rightarrow b) = (a - b) * m$$

<proof>

lemma *iMODb-inext-nth-diff*: $\llbracket a \leq c; b \leq c \rrbracket \implies$

$$([r, \text{mod } m, c] \rightarrow a) - ([r, \text{mod } m, c] \rightarrow b) = (a - b) * m$$

<proof>

1.8 Induction with intervals

lemma *iFROM-induct*:

$$\llbracket P n; \bigwedge k. \llbracket k \in [n..]; P k \rrbracket \implies P (\text{Suc } k); a \in [n..] \rrbracket \implies P a$$

<proof>

lemma *iIN-induct*:

$$\llbracket P n; \bigwedge k. \llbracket k \in [n.., d]; k \neq n + d; P k \rrbracket \implies P (\text{Suc } k); a \in [n.., d] \rrbracket \implies P a$$

<proof>

lemma *iTILL-induct*:

$$\llbracket P 0; \bigwedge k. \llbracket k \in [..n]; k \neq n; P k \rrbracket \implies P (\text{Suc } k); a \in [..n] \rrbracket \implies P a$$

<proof>

lemma *iMOD-induct*:

$$\llbracket P r; \bigwedge k. \llbracket k \in [r, \text{mod } m]; P k \rrbracket \implies P (k + m); a \in [r, \text{mod } m] \rrbracket \implies P a$$

<proof>

lemma *iMODb-induct*:

$$\llbracket P r; \bigwedge k. \llbracket k \in [r, \text{mod } m, c]; k \neq r + m * c; P k \rrbracket \implies P (k + m); a \in [r, \text{mod } m, c] \rrbracket \implies P a$$

<proof>

lemma *iIN-rev-induct*:

$$\llbracket P (n + d); \bigwedge k. \llbracket k \in [n.., d]; k \neq n; P k \rrbracket \implies P (k - \text{Suc } 0); a \in [n.., d] \rrbracket \implies P a$$

<proof>

lemma *iTILL-rev-induct*:

$$\llbracket P n; \bigwedge k. \llbracket k \in [..n]; 0 < k; P k \rrbracket \implies P (k - \text{Suc } 0); a \in [..n] \rrbracket \implies P a$$

<proof>

lemma *iMODb-rev-induct*:

$$\llbracket P (r + m * c); \bigwedge k. \llbracket k \in [r, \text{mod } m, c]; k \neq r; P k \rrbracket \implies P (k - m); a \in [r, \text{mod } m, c] \rrbracket \implies P a$$
 $\langle \text{proof} \rangle$

end

2 Arithmetic operators on natural intervals

theory *IL-IntervalOperators*
imports *IL-Interval*
begin

2.1 Arithmetic operations with intervals

2.1.1 Addition of and multiplication by constants

definition *iT-Plus* :: $iT \Rightarrow Time \Rightarrow iT$ (**infixl** $\langle \oplus \rangle$ 55)
where $I \oplus k \equiv (\lambda n. (n + k)) \text{ ' } I$

definition *iT-Mult* :: $iT \Rightarrow Time \Rightarrow iT$ (**infixl** $\langle \otimes \rangle$ 55)
where $iT\text{-Mult-def} : I \otimes k \equiv (\lambda n. (n * k)) \text{ ' } I$

lemma *iT-Plus-image-conv*: $I \oplus k = (\lambda n. (n + k)) \text{ ' } I$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-image-conv*: $I \otimes k = (\lambda n. (n * k)) \text{ ' } I$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-empty*: $\{\} \oplus k = \{\}$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-empty*: $\{\} \otimes k = \{\}$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-not-empty*: $I \neq \{\} \implies I \oplus k \neq \{\}$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-not-empty*: $I \neq \{\} \implies I \otimes k \neq \{\}$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-empty-iff*: $(I \oplus k = \{\}) = (I = \{\})$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-empty-iff*: $(I \otimes k = \{\}) = (I = \{\})$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-mono*: $A \subseteq B \implies A \oplus k \subseteq B \oplus k$
 $\langle proof \rangle$

lemma *iT-Mult-mono*: $A \subseteq B \implies A \otimes k \subseteq B \otimes k$
 $\langle proof \rangle$

lemma *iT-Mult-0*: $I \neq \{\} \implies I \otimes 0 = [\dots 0]$
 $\langle proof \rangle$

corollary *iT-Mult-0-if*: $I \otimes 0 = (\text{if } I = \{\} \text{ then } \{\} \text{ else } [\dots 0])$
 $\langle proof \rangle$

lemma *iT-Plus-mem-iff*: $x \in (I \oplus k) = (k \leq x \wedge (x - k) \in I)$
 $\langle proof \rangle$

lemma *iT-Plus-mem-iff2*: $x + k \in (I \oplus k) = (x \in I)$
 $\langle proof \rangle$

lemma *iT-Mult-mem-iff-0*: $x \in (I \otimes 0) = (I \neq \{\} \wedge x = 0)$
 $\langle proof \rangle$

lemma *iT-Mult-mem-iff*:
 $0 < k \implies x \in (I \otimes k) = (x \text{ mod } k = 0 \wedge x \text{ div } k \in I)$
 $\langle proof \rangle$

lemma *iT-Mult-mem-iff2*: $0 < k \implies x * k \in (I \otimes k) = (x \in I)$
 $\langle proof \rangle$

lemma *iT-Plus-singleton*: $\{a\} \oplus k = \{a + k\}$
 $\langle proof \rangle$

lemma *iT-Mult-singleton*: $\{a\} \otimes k = \{a * k\}$
 $\langle proof \rangle$

lemma *iT-Plus-Un*: $(A \cup B) \oplus k = (A \oplus k) \cup (B \oplus k)$
 $\langle proof \rangle$

lemma *iT-Mult-Un*: $(A \cup B) \otimes k = (A \otimes k) \cup (B \otimes k)$
 $\langle proof \rangle$

lemma *iT-Plus-Int*: $(A \cap B) \oplus k = (A \oplus k) \cap (B \oplus k)$
 $\langle proof \rangle$

lemma *iT-Mult-Int*: $0 < k \implies (A \cap B) \otimes k = (A \otimes k) \cap (B \otimes k)$
 $\langle proof \rangle$

lemma *iT-Plus-image*: $f \text{ ‘ } I \oplus n = (\lambda x. f\ x + n) \text{ ‘ } I$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-image*: $f \text{ ‘ } I \otimes n = (\lambda x. f\ x * n) \text{ ‘ } I$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-commute*: $I \oplus a \oplus b = I \oplus b \oplus a$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-commute*: $I \otimes a \otimes b = I \otimes b \otimes a$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-assoc*: $I \oplus a \oplus b = I \oplus (a + b)$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-assoc*: $I \otimes a \otimes b = I \otimes (a * b)$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-Mult-distrib*: $I \oplus n \otimes m = I \otimes m \oplus n * m$
 $\langle \text{proof} \rangle \langle \text{proof} \rangle \langle \text{proof} \rangle \langle \text{proof} \rangle \langle \text{proof} \rangle \langle \text{proof} \rangle$

lemma *iT-Plus-finite-iff*: $\text{finite } (I \oplus k) = \text{finite } I$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-0-finite*: $\text{finite } (I \otimes 0)$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-finite-iff*: $0 < k \implies \text{finite } (I \otimes k) = \text{finite } I$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-Min*: $I \neq \{\} \implies iMin\ (I \oplus k) = iMin\ I + k$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-Min*: $I \neq \{\} \implies iMin\ (I \otimes k) = iMin\ I * k$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-Max*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies Max\ (I \oplus k) = Max\ I + k$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-Max*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies Max\ (I \otimes k) = Max\ I * k$
 $\langle \text{proof} \rangle$

corollary

iMOD-mult-0: $[r, \text{mod } m] \otimes 0 = [\dots 0]$ **and**
iMODb-mult-0: $[r, \text{mod } m, c] \otimes 0 = [\dots 0]$ **and**
iFROM-mult-0: $[n \dots] \otimes 0 = [\dots 0]$ **and**
iIN-mult-0: $[n \dots, d] \otimes 0 = [\dots 0]$ **and**
iTILL-mult-0: $[\dots n] \otimes 0 = [\dots 0]$

$\langle \text{proof} \rangle$

lemmas *iT-mult-0* =
iTILL-mult-0
iFROM-mult-0
iIN-mult-0
iMOD-mult-0
iMODb-mult-0

lemma *iT-Plus-0*: $I \oplus 0 = I$
 $\langle proof \rangle$

lemma *iT-Mult-1*: $I \otimes Suc\ 0 = I$
 $\langle proof \rangle$

corollary
iFROM-add-Min: $iMin\ ([n..] \oplus k) = n + k$ **and**
iIN-add-Min: $iMin\ ([n..,d] \oplus k) = n + k$ **and**
iTILL-add-Min: $iMin\ ([..n] \oplus k) = k$ **and**
iMOD-add-Min: $iMin\ ([r, mod\ m] \oplus k) = r + k$ **and**
iMODb-add-Min: $iMin\ ([r, mod\ m, c] \oplus k) = r + k$
 $\langle proof \rangle$

corollary
iFROM-mult-Min: $iMin\ ([n..] \otimes k) = n * k$ **and**
iIN-mult-Min: $iMin\ ([n..,d] \otimes k) = n * k$ **and**
iTILL-mult-Min: $iMin\ ([..n] \otimes k) = 0$ **and**
iMOD-mult-Min: $iMin\ ([r, mod\ m] \otimes k) = r * k$ **and**
iMODb-mult-Min: $iMin\ ([r, mod\ m, c] \otimes k) = r * k$
 $\langle proof \rangle$

lemmas *iT-add-Min* =
iIN-add-Min
iTILL-add-Min
iFROM-add-Min
iMOD-add-Min
iMODb-add-Min

lemmas *iT-mult-Min* =
iIN-mult-Min
iTILL-mult-Min
iFROM-mult-Min
iMOD-mult-Min
iMODb-mult-Min

lemma *iFROM-add*: $[n..] \oplus k = [n+k..]$
 $\langle proof \rangle$

lemma *iIN-add*: $[n..,d] \oplus k = [n+k..,d]$

$\langle proof \rangle$

lemma *iTILL-add*: $[\dots i] \oplus k = [k\dots, i]$
 $\langle proof \rangle$

lemma *iMOD-add*: $[r, \text{mod } m] \oplus k = [r + k, \text{mod } m]$
 $\langle proof \rangle$

lemma *iMODb-add*: $[r, \text{mod } m, c] \oplus k = [r + k, \text{mod } m, c]$
 $\langle proof \rangle$

lemmas *iT-add* =
iMOD-add
iMODb-add
iFROM-add
iIN-add
iTILL-add
iT-Plus-singleton

lemma *iFROM-mult*: $[n\dots] \otimes k = [n * k, \text{mod } k]$
 $\langle proof \rangle$

lemma *iIN-mult*: $[n\dots, d] \otimes k = [n * k, \text{mod } k, d]$
 $\langle proof \rangle$

lemma *iTILL-mult*: $[\dots n] \otimes k = [0, \text{mod } k, n]$
 $\langle proof \rangle$

lemma *iMOD-mult*: $[r, \text{mod } m] \otimes k = [r * k, \text{mod } m * k]$
 $\langle proof \rangle$

lemma *iMODb-mult*:
 $[r, \text{mod } m, c] \otimes k = [r * k, \text{mod } m * k, c]$
 $\langle proof \rangle$

lemmas *iT-mult* =
iTILL-mult
iFROM-mult
iIN-mult
iMOD-mult
iMODb-mult
iT-Mult-singleton

2.1.2 Some conversions between intervals using constant addition and multiplication

lemma *iFROM-conv*: $[n\dots] = UNIV \oplus n$
 $\langle proof \rangle$

lemma *iN-conv*: $[n \dots d] = [\dots d] \oplus n$
 $\langle \text{proof} \rangle$

lemma *iMOD-conv*: $[r, \text{mod } m] = [0 \dots] \otimes m \oplus r$
 $\langle \text{proof} \rangle$

lemma *iMODb-conv*: $[r, \text{mod } m, c] = [\dots c] \otimes m \oplus r$
 $\langle \text{proof} \rangle$

Some examples showing the utility of iMODb_conv

lemma $[12, \text{mod } 10, 4] = \{12, 22, 32, 42, 52\}$
 $\langle \text{proof} \rangle$

lemma $[12, \text{mod } 10, 4] = \{12, 22, 32, 42, 52\}$
 $\langle \text{proof} \rangle$

lemma $[12, \text{mod } 10, 4] = \{12, 22, 32, 42, 52\}$
 $\langle \text{proof} \rangle$

lemma $[r, \text{mod } m, 4] = \{r, r+m, r+2*m, r+3*m, r+4*m\}$
 $\langle \text{proof} \rangle$

lemma $[2, \text{mod } 10, 4] = \{2, 12, 22, 32, 42\}$
 $\langle \text{proof} \rangle$

2.1.3 Subtraction of constants

definition *iT-Plus-neg* :: $iT \Rightarrow \text{Time} \Rightarrow iT$ (**infixl** $\langle \oplus - \rangle$ 55) **where**
 $I \oplus - k \equiv \{x. x + k \in I\}$

lemma *iT-Plus-neg-mem-iff*: $(x \in I \oplus - k) = (x + k \in I)$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-mem-iff2*: $k \leq x \implies (x - k \in I \oplus - k) = (x \in I)$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-image-conv*: $I \oplus - k = (\lambda n. (n - k)) \text{ ' } (I \downarrow \geq k)$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-cut-eq*: $t \leq k \implies (I \downarrow \geq t) \oplus - k = I \oplus - k$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-mono*: $A \subseteq B \implies A \oplus - k \subseteq B \oplus - k$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-empty*: $\{\} \oplus - k = \{\}$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-Max-less-empty*:

$$\llbracket \text{finite } I; \text{Max } I < k \rrbracket \implies I \oplus - k = \{\}$$

<proof>

lemma *iT-Plus-neg-not-empty-iff*: $(I \oplus - k \neq \{\}) = (\exists x \in I. k \leq x)$

<proof>

lemma *iT-Plus-neg-empty-iff*:

$$(I \oplus - k = \{\}) = (I = \{\} \vee (\text{finite } I \wedge \text{Max } I < k))$$

<proof>

lemma *iT-Plus-neg-assoc*: $(I \oplus - a) \oplus - b = I \oplus - (a + b)$

<proof>

lemma *iT-Plus-neg-commute*: $I \oplus - a \oplus - b = I \oplus - b \oplus - a$

<proof>

lemma *iT-Plus-neg-0*: $I \oplus - 0 = I$

<proof>

lemma *iT-Plus-Plus-neg-assoc*: $b \leq a \implies I \oplus a \oplus - b = I \oplus (a - b)$

<proof>

lemma *iT-Plus-Plus-neg-assoc2*: $a \leq b \implies I \oplus a \oplus - b = I \oplus - (b - a)$

<proof>

lemma *iT-Plus-neg-Plus-le-cut-eq*:

$$a \leq b \implies (I \oplus - a) \oplus b = (I \downarrow \geq a) \oplus (b - a)$$

<proof>

corollary *iT-Plus-neg-Plus-le-Min-eq*:

$$\llbracket a \leq b; a \leq \text{iMin } I \rrbracket \implies (I \oplus - a) \oplus b = I \oplus (b - a)$$

<proof>

lemma *iT-Plus-neg-Plus-ge-cut-eq*:

$$b \leq a \implies (I \oplus - a) \oplus b = (I \downarrow \geq a) \oplus - (a - b)$$

<proof>

corollary *iT-Plus-neg-Plus-ge-Min-eq*:

$$\llbracket b \leq a; a \leq \text{iMin } I \rrbracket \implies (I \oplus - a) \oplus b = I \oplus - (a - b)$$

<proof>

lemma *iT-Plus-neg-Mult-distrib*:

$$0 < m \implies I \oplus - n \otimes m = I \otimes m \oplus - n * m$$

<proof>

lemma *iT-Plus-neg-Plus-le-inverse*: $k \leq \text{iMin } I \implies I \oplus - k \oplus k = I$

<proof>

lemma *iT-Plus-neg-Plus-inverse*: $I \oplus - k \oplus k = I \downarrow \geq k$

$\langle proof \rangle$

lemma *iT-Plus-Plus-neg-inverse*: $I \oplus k \oplus - k = I$
 $\langle proof \rangle$

lemma *iT-Plus-neg-Un*: $(A \cup B) \oplus - k = (A \oplus - k) \cup (B \oplus - k)$
 $\langle proof \rangle$

lemma *iT-Plus-neg-Int*: $(A \cap B) \oplus - k = (A \oplus - k) \cap (B \oplus - k)$
 $\langle proof \rangle$

lemma *iT-Plus-neg-Max-singleton*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies I \oplus - \text{Max } I = \{0\}$
 $\langle proof \rangle$

lemma *iT-Plus-neg-singleton*: $\{a\} \oplus - k = (\text{if } k \leq a \text{ then } \{a - k\} \text{ else } \{\})$
 $\langle proof \rangle$

corollary *iT-Plus-neg-singleton1*: $k \leq a \implies \{a\} \oplus - k = \{a - k\}$
 $\langle proof \rangle$

corollary *iT-Plus-neg-singleton2*: $a < k \implies \{a\} \oplus - k = \{\}$
 $\langle proof \rangle$

lemma *iT-Plus-neg-finite-iff*: $\text{finite } (I \oplus - k) = \text{finite } I$
 $\langle proof \rangle$

lemma *iT-Plus-neg-Min*:
 $I \oplus - k \neq \{\} \implies iMin (I \oplus - k) = iMin (I \downarrow \geq k) - k$
 $\langle proof \rangle$

lemma *iT-Plus-neg-Max*:
 $\llbracket \text{finite } I; I \oplus - k \neq \{\} \rrbracket \implies \text{Max } (I \oplus - k) = \text{Max } I - k$
 $\langle proof \rangle$

Subtractions of constants from intervals

lemma *iFROM-add-neg*: $[n \dots] \oplus - k = [n - k \dots]$
 $\langle proof \rangle$

lemma *iTILL-add-neg*: $[\dots n] \oplus - k = (\text{if } k \leq n \text{ then } [\dots n - k] \text{ else } \{\})$
 $\langle proof \rangle$

lemma *iTILL-add-neg1*: $k \leq n \implies [\dots n] \oplus - k = [\dots n - k]$
 $\langle proof \rangle$

lemma *iTILL-add-neg2*: $n < k \implies [\dots n] \oplus - k = \{\}$
 $\langle proof \rangle$

lemma *iIN-add-neg*:
 $[n \dots, d] \oplus - k = ($
 $\text{if } k \leq n \text{ then } [n - k \dots, d]$

else if $k \leq n + d$ then $[\dots n + d - k]$ else $\{\}$
 $\langle \text{proof} \rangle$

lemma *iIN-add-neg1*: $k \leq n \implies [n\dots, d] \oplus -k = [n - k\dots, d]$
 $\langle \text{proof} \rangle$

lemma *iIN-add-neg2*: $\llbracket n \leq k; k \leq n + d \rrbracket \implies [n\dots, d] \oplus -k = [\dots n + d - k]$
 $\langle \text{proof} \rangle$

lemma *iIN-add-neg3*: $n + d < k \implies [n\dots, d] \oplus -k = \{\}$
 $\langle \text{proof} \rangle$

lemma *iMOD-0-add-neg*: $[r, \text{mod } 0] \oplus -k = \{r\} \oplus -k$
 $\langle \text{proof} \rangle$

lemma *iMOD-gr0-add-neg*:
 $0 < m \implies$
 $[r, \text{mod } m] \oplus -k = ($
 $\quad \text{if } k \leq r \text{ then } [r - k, \text{mod } m]$
 $\quad \text{else } [(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m, \text{mod } m])$
 $\langle \text{proof} \rangle$

lemma *iMOD-add-neg*:
 $[r, \text{mod } m] \oplus -k = ($
 $\quad \text{if } k \leq r \text{ then } [r - k, \text{mod } m]$
 $\quad \text{else if } 0 < m \text{ then } [(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m, \text{mod } m] \text{ else } \{\})$
 $\langle \text{proof} \rangle$

corollary *iMOD-add-neg1*:
 $k \leq r \implies [r, \text{mod } m] \oplus -k = [r - k, \text{mod } m]$
 $\langle \text{proof} \rangle$

lemma *iMOD-add-neg2*:
 $\llbracket 0 < m; r < k \rrbracket \implies [r, \text{mod } m] \oplus -k = [(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m,$
 $\text{mod } m]$
 $\langle \text{proof} \rangle$

lemma *iMODb-mod-0-add-neg*: $[r, \text{mod } 0, c] \oplus -k = \{r\} \oplus -k$
 $\langle \text{proof} \rangle$

lemma *iMODb-add-neg*:
 $[r, \text{mod } m, c] \oplus -k = ($
 $\quad \text{if } k \leq r \text{ then } [r - k, \text{mod } m, c]$
 $\quad \text{else}$

$$\begin{aligned} & \text{if } k \leq r + m * c \text{ then} \\ & \quad [(m + r \bmod m - k \bmod m) \bmod m, \bmod m, (r + m * c - k) \operatorname{div} m] \\ & \text{else } \{\} \end{aligned}$$
 $\langle \text{proof} \rangle$

lemma *iMODb-add-neg'*:

$$\begin{aligned} & [r, \bmod m, c] \oplus -k = (\\ & \quad \text{if } k \leq r \text{ then } [r - k, \bmod m, c] \\ & \quad \text{else if } k \leq r + m * c \text{ then} \\ & \quad \quad \text{if } k \bmod m \leq r \bmod m \\ & \quad \quad \quad \text{then } [r \bmod m - k \bmod m, \bmod m, c + r \operatorname{div} m - k \operatorname{div} m] \\ & \quad \quad \quad \text{else } [m + r \bmod m - k \bmod m, \bmod m, c + r \operatorname{div} m - \operatorname{Suc}(k \operatorname{div} m)] \\ & \quad \text{else } \{\} \end{aligned}$$
 $\langle \text{proof} \rangle$

corollary *iMODb-add-neg1*:

$$k \leq r \implies [r, \bmod m, c] \oplus -k = [r - k, \bmod m, c]$$
 $\langle \text{proof} \rangle$

corollary *iMODb-add-neg2*:

$$\begin{aligned} & \llbracket r < k; k \leq r + m * c \rrbracket \implies \\ & [r, \bmod m, c] \oplus -k = \\ & [(m + r \bmod m - k \bmod m) \bmod m, \bmod m, (r + m * c - k) \operatorname{div} m] \end{aligned}$$
 $\langle \text{proof} \rangle$

corollary *iMODb-add-neg2-mod-le*:

$$\begin{aligned} & \llbracket r < k; k \leq r + m * c; k \bmod m \leq r \bmod m \rrbracket \implies \\ & [r, \bmod m, c] \oplus -k = \\ & [r \bmod m - k \bmod m, \bmod m, c + r \operatorname{div} m - k \operatorname{div} m] \end{aligned}$$
 $\langle \text{proof} \rangle$

corollary *iMODb-add-neg2-mod-less*:

$$\begin{aligned} & \llbracket r < k; k \leq r + m * c; r \bmod m < k \bmod m \rrbracket \implies \\ & [r, \bmod m, c] \oplus -k = \\ & [m + r \bmod m - k \bmod m, \bmod m, c + r \operatorname{div} m - \operatorname{Suc}(k \operatorname{div} m)] \end{aligned}$$
 $\langle \text{proof} \rangle$

lemma *iMODb-add-neg3*: $r + m * c < k \implies [r, \bmod m, c] \oplus -k = \{\}$

$\langle \text{proof} \rangle$

lemmas *iT-add-neg* =

iFROM-add-neg
iIN-add-neg
iTILL-add-neg
iMOD-add-neg
iMODb-add-neg
iT-Plus-neg-singleton

2.1.4 Subtraction of intervals from constants

definition $iT\text{-Minus} :: Time \Rightarrow iT \Rightarrow iT$ (**infixl** $\ominus$ 55)
where $k \ominus I \equiv \{x. x \leq k \wedge (k - x) \in I\}$

lemma $iT\text{-Minus-mem-iff}$: $(x \in k \ominus I) = (x \leq k \wedge k - x \in I)$
 $\langle proof \rangle$

lemma $iT\text{-Minus-mono}$: $A \subseteq B \implies k \ominus A \subseteq k \ominus B$
 $\langle proof \rangle$

lemma $iT\text{-Minus-image-conv}$: $k \ominus I = (\lambda x. k - x) ` (I \downarrow \leq k)$
 $\langle proof \rangle$

lemma $iT\text{-Minus-cut-eq}$: $k \leq t \implies k \ominus (I \downarrow \leq t) = k \ominus I$
 $\langle proof \rangle$

lemma $iT\text{-Minus-Minus-cut-eq}$: $k \ominus (k \ominus (I \downarrow \leq k)) = I \downarrow \leq k$
 $\langle proof \rangle$

lemma $10 \ominus [\dots 3] = [7 \dots, 3]$
 $\langle proof \rangle$

lemma $iT\text{-Minus-empty}$: $k \ominus \{\} = \{\}$
 $\langle proof \rangle$

lemma $iT\text{-Minus-0}$: $k \ominus \{0\} = \{k\}$
 $\langle proof \rangle$

lemma $iT\text{-Minus-bound}$: $x \in k \ominus I \implies x \leq k$
 $\langle proof \rangle$

lemma $iT\text{-Minus-finite}$: $finite (k \ominus I)$
 $\langle proof \rangle$

lemma $iT\text{-Minus-less-Min-empty}$: $k < iMin I \implies k \ominus I = \{\}$
 $\langle proof \rangle$

lemma $iT\text{-Minus-Min-singleton}$: $I \neq \{\} \implies (iMin I) \ominus I = \{0\}$
 $\langle proof \rangle$

lemma $iT\text{-Minus-empty-iff}$: $(k \ominus I = \{\}) = (I = \{\} \vee k < iMin I)$
 $\langle proof \rangle$

lemma $iT\text{-Minus-imirror-conv}$:
 $k \ominus I = imirror (I \downarrow \leq k) \oplus k \oplus - (iMin I + Max (I \downarrow \leq k))$
 $\langle proof \rangle$

lemma *iT-Minus-imirror-conv'*:

$$k \ominus I = \text{imirror } (I \downarrow \leq k) \oplus k \oplus - (iMin (I \downarrow \leq k) + Max (I \downarrow \leq k))$$

<proof>

lemma *iT-Minus-Max*:

$$\llbracket I \neq \{\}; iMin I \leq k \rrbracket \implies Max (k \ominus I) = k - (iMin I)$$

<proof>

lemma *iT-Minus-Min*:

$$\llbracket I \neq \{\}; iMin I \leq k \rrbracket \implies iMin (k \ominus I) = k - (Max (I \downarrow \leq k))$$

<proof>

lemma *iT-Minus-Minus-eq*: $\llbracket \text{finite } I; Max I \leq k \rrbracket \implies k \ominus (k \ominus I) = I$

<proof>

lemma *iT-Minus-Minus-eq2*: $I \subseteq [\dots k] \implies k \ominus (k \ominus I) = I$

<proof>

lemma *iT-Minus-Minus*: $a \ominus (b \ominus I) = (I \downarrow \leq b) \oplus a \oplus - b$

<proof>

lemma *iT-Minus-Plus-empty*: $k < n \implies k \ominus (I \oplus n) = \{\}$

<proof>

lemma *iT-Minus-Plus-commute*: $n \leq k \implies k \ominus (I \oplus n) = (k - n) \ominus I$

<proof>

lemma *iT-Minus-Plus-cut-assoc*: $(k \ominus I) \oplus n = (k + n) \ominus (I \downarrow \leq k)$

<proof>

lemma *iT-Minus-Plus-assoc*:

$$\llbracket \text{finite } I; Max I \leq k \rrbracket \implies (k \ominus I) \oplus n = (k + n) \ominus I$$

<proof>

lemma *iT-Minus-Plus-assoc2*:

$$I \subseteq [\dots k] \implies (k \ominus I) \oplus n = (k + n) \ominus I$$

<proof>

lemma *iT-Minus-Un*: $k \ominus (A \cup B) = (k \ominus A) \cup (k \ominus B)$

<proof>

lemma *iT-Minus-Int*: $k \ominus (A \cap B) = (k \ominus A) \cap (k \ominus B)$

<proof>

lemma *iT-Minus-singleton*: $k \ominus \{a\} = (\text{if } a \leq k \text{ then } \{k - a\} \text{ else } \{\})$

<proof>

corollary *iT-Minus-singleton1*: $a \leq k \implies k \ominus \{a\} = \{k - a\}$

<proof>

corollary *iT-Minus-singleton2*: $k < a \implies k \ominus \{a\} = \{\}$

$\langle proof \rangle$

lemma *iMOD-sub*:

$k \ominus [r, \text{mod } m] =$
 $(\text{if } r \leq k \text{ then } [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m] \text{ else } \{\})$
 $\langle proof \rangle$

corollary *iMOD-sub1*:

$r \leq k \implies k \ominus [r, \text{mod } m] = [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m]$
 $\langle proof \rangle$

corollary *iMOD-sub2*: $k < r \implies k \ominus [r, \text{mod } m] = \{\}$
 $\langle proof \rangle$

lemma *iTILL-sub*: $k \ominus [\dots n] = (\text{if } n \leq k \text{ then } [k - n \dots, n] \text{ else } [\dots k])$
 $\langle proof \rangle$

corollary *iTILL-sub1*: $n \leq k \implies k \ominus [\dots n] = [k - n \dots, n]$
 $\langle proof \rangle$

corollary *iTILL-sub2*: $k \leq n \implies k \ominus [\dots n] = [\dots k]$
 $\langle proof \rangle$

lemma *iMODb-sub*:

$k \ominus [r, \text{mod } m, c] =$
 $(\text{if } r + m * c \leq k \text{ then } [k - r - m * c, \text{mod } m, c] \text{ else}$
 $\text{if } r \leq k \text{ then } [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m] \text{ else } \{\})$
 $\langle proof \rangle$

corollary *iMODb-sub1*:

$\llbracket r \leq k; k \leq r + m * c \rrbracket \implies$
 $k \ominus [r, \text{mod } m, c] = [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m]$
 $\langle proof \rangle$

corollary *iMODb-sub2*: $k < r \implies k \ominus [r, \text{mod } m, c] = \{\}$
 $\langle proof \rangle$

corollary *iMODb-sub3*:

$r + m * c \leq k \implies k \ominus [r, \text{mod } m, c] = [k - r - m * c, \text{mod } m, c]$
 $\langle proof \rangle$

lemma *iFROM-sub*: $k \ominus [n \dots] = (\text{if } n \leq k \text{ then } [\dots k - n] \text{ else } \{\})$
 $\langle proof \rangle$

corollary *iFROM-sub1*: $n \leq k \implies k \ominus [n \dots] = [\dots k - n]$
 $\langle proof \rangle$

corollary *iFROM-sub-empty*: $k < n \implies k \ominus [n..] = \{\}$
 $\langle \text{proof} \rangle$

lemma *iIN-sub*:
 $k \ominus [n.., d] = ($
 $\text{if } n + d \leq k \text{ then } [k - (n + d).., d]$
 $\text{else if } n \leq k \text{ then } [\dots k - n] \text{ else } \{\})$
 $\langle \text{proof} \rangle$

lemma *iIN-sub1*: $n + d \leq k \implies k \ominus [n.., d] = [k - (n + d).., d]$
 $\langle \text{proof} \rangle$

lemma *iIN-sub2*: $\llbracket n \leq k; k \leq n + d \rrbracket \implies k \ominus [n.., d] = [\dots k - n]$
 $\langle \text{proof} \rangle$

lemma *iIN-sub3*: $k < n \implies k \ominus [n.., d] = \{\}$
 $\langle \text{proof} \rangle$

lemmas *iT-sub* =
iFROM-sub
iIN-sub
iTILL-sub
iMOD-sub
iMODb-sub
iT-Minus-singleton

2.1.5 Division of intervals by constants

Monotonicity and injectivity of arithmetic operators

lemma *iMOD-div-right-strict-mono-on*:
 $\llbracket 0 < k; k \leq m \rrbracket \implies \text{strict-mono-on } (\lambda x. x \text{ div } k) [r, \text{ mod } m]$
 $\langle \text{proof} \rangle$

corollary *iMOD-div-right-inj-on*:
 $\llbracket 0 < k; k \leq m \rrbracket \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{ mod } m]$
 $\langle \text{proof} \rangle$

lemma *iMOD-mult-div-right-inj-on*:
 $\text{inj-on } (\lambda x. x \text{ div } (k::\text{nat})) [r, \text{ mod } (k * m)]$
 $\langle \text{proof} \rangle$

lemma *iMOD-mult-div-right-inj-on2*:
 $m \text{ mod } k = 0 \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{ mod } m]$
 $\langle \text{proof} \rangle$

lemma *iMODb-div-right-strict-mono-on*:

$\llbracket 0 < k; k \leq m \rrbracket \implies \text{strict-mono-on } (\lambda x. x \text{ div } k) [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

corollary *iMODb-div-right-inj-on*:

$\llbracket 0 < k; k \leq m \rrbracket \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

lemma *iMODb-mult-div-right-inj-on*:

$\text{inj-on } (\lambda x. x \text{ div } (k::\text{nat})) [r, \text{mod } (k * m), c]$
 $\langle \text{proof} \rangle$

corollary *iMODb-mult-div-right-inj-on2*:

$m \text{ mod } k = 0 \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

definition *iT-Div* :: *iT* $\Rightarrow$ *Time* $\Rightarrow$ *iT* (**infixl** $\circledast$ 55)

where $I \circledast k \equiv (\lambda n. (n \text{ div } k)) \circ I$

lemma *iT-Div-image-conv*: $I \circledast k = (\lambda n. (n \text{ div } k)) \circ I$

$\langle \text{proof} \rangle$

lemma *iT-Div-mono*: $A \subseteq B \implies A \circledast k \subseteq B \circledast k$

$\langle \text{proof} \rangle$

lemma *iT-Div-empty*: $\{\} \circledast k = \{\}$

$\langle \text{proof} \rangle$

lemma *iT-Div-not-empty*: $I \neq \{\} \implies I \circledast k \neq \{\}$

$\langle \text{proof} \rangle$

lemma *iT-Div-empty-iff*: $(I \circledast k = \{\}) = (I = \{\})$

$\langle \text{proof} \rangle$

lemma *iT-Div-0*: $I \neq \{\} \implies I \circledast 0 = [\dots 0]$

$\langle \text{proof} \rangle$

corollary *iT-Div-0-if*: $I \circledast 0 = (\text{if } I = \{\} \text{ then } \{\} \text{ else } [\dots 0])$

$\langle \text{proof} \rangle$

corollary

iFROM-div-0: $[n\dots] \circledast 0 = [\dots 0]$ **and**

iTILL-div-0: $[\dots n] \circledast 0 = [\dots 0]$ **and**

iIN-div-0: $[n\dots, d] \circledast 0 = [\dots 0]$ **and**

iMOD-div-0: $[r, \text{mod } m] \circledast 0 = [\dots 0]$ **and**

iMODb-div-0: $[r, \text{mod } m, c] \circledast 0 = [\dots 0]$

$\langle \text{proof} \rangle$

lemmas *iT-div-0* =

iTILL-div-0

iFROM-div-0

iIN-div-0
iMOD-div-0
iMODb-div-0

lemma *iT-Div-1*: $I \oslash \text{Suc } 0 = I$
 $\langle \text{proof} \rangle$

lemma *iT-Div-mem-iff-0*: $x \in (I \oslash 0) = (I \neq \{\} \wedge x = 0)$
 $\langle \text{proof} \rangle$

lemma *iT-Div-mem-iff*:
 $0 < k \implies x \in (I \oslash k) = (\exists y \in I. y \text{ div } k = x)$
 $\langle \text{proof} \rangle$

lemma *iT-Div-mem-iff2*:
 $0 < k \implies x \text{ div } k \in (I \oslash k) = (\exists y \in I. y \text{ div } k = x \text{ div } k)$
 $\langle \text{proof} \rangle$

lemma *iT-Div-mem-iff-Int*:
 $0 < k \implies x \in (I \oslash k) = (I \cap [x * k \dots k - \text{Suc } 0] \neq \{\})$
 $\langle \text{proof} \rangle$

lemma *iT-Div-imp-mem*:
 $0 < k \implies x \in I \implies x \text{ div } k \in (I \oslash k)$
 $\langle \text{proof} \rangle$

lemma *iT-Div-singleton*: $\{a\} \oslash k = \{a \text{ div } k\}$
 $\langle \text{proof} \rangle$

lemma *iT-Div-Un*: $(A \cup B) \oslash k = (A \oslash k) \cup (B \oslash k)$
 $\langle \text{proof} \rangle$

lemma *iT-Div-insert*: $(\text{insert } n \ I) \oslash k = \text{insert } (n \text{ div } k) \ (I \oslash k)$
 $\langle \text{proof} \rangle$

lemma *not-iT-Div-Int*: $\neg (\forall k \ A \ B. (A \cap B) \oslash k = (A \oslash k) \cap (B \oslash k))$
 $\langle \text{proof} \rangle$

lemma *subset-iT-Div-Int*: $A \subseteq B \implies (A \cap B) \oslash k = (A \oslash k) \cap (B \oslash k)$
 $\langle \text{proof} \rangle$

lemma *iFROM-iT-Div-Int*:
 $\llbracket 0 < k; n \leq iMin \ A \rrbracket \implies (A \cap [n \dots]) \oslash k = (A \oslash k) \cap ([n \dots] \oslash k)$
 $\langle \text{proof} \rangle$

lemma *iIN-iT-Div-Int*:
 $\llbracket 0 < k; n \leq iMin \ A; \forall x \in A. x \text{ div } k \leq (n + d) \text{ div } k \longrightarrow x \leq n + d \rrbracket \implies$

$$(A \cap [n..,d]) \odot k = (A \odot k) \cap ([n..,d] \odot k)$$

<proof>

corollary *iTILL-iT-Div-Int:*

$$\llbracket 0 < k; \forall x \in A. x \text{ div } k \leq n \text{ div } k \longrightarrow x \leq n \rrbracket \Longrightarrow$$

$$(A \cap [\dots n]) \odot k = (A \odot k) \cap ([\dots n] \odot k)$$

<proof>

lemma *iIN-iT-Div-Int-mod-0:*

$$\llbracket 0 < k; n \bmod k = 0; \forall x \in A. x \text{ div } k \leq (n + d) \text{ div } k \longrightarrow x \leq n + d \rrbracket \Longrightarrow$$

$$(A \cap [n..,d]) \odot k = (A \odot k) \cap ([n..,d] \odot k)$$

<proof>

lemma *mod-partition-iT-Div-Int:*

$$\llbracket 0 < k; 0 < d \rrbracket \Longrightarrow$$

$$(A \cap [n * k.., d * k - \text{Suc } 0]) \odot k =$$

$$(A \odot k) \cap ([n * k.., d * k - \text{Suc } 0] \odot k)$$

<proof> *<proof>*

corollary *mod-partition-iT-Div-Int2:*

$$\llbracket 0 < k; 0 < d; n \bmod k = 0; d \bmod k = 0 \rrbracket \Longrightarrow$$

$$(A \cap [n.., d - \text{Suc } 0]) \odot k =$$

$$(A \odot k) \cap ([n.., d - \text{Suc } 0] \odot k)$$

<proof>

corollary *mod-partition-iT-Div-Int-one-segment:*

$$0 < k \Longrightarrow$$

$$(A \cap [n * k.., k - \text{Suc } 0]) \odot k = (A \odot k) \cap ([n * k.., k - \text{Suc } 0] \odot k)$$

<proof>

corollary *mod-partition-iT-Div-Int-one-segment2:*

$$\llbracket 0 < k; n \bmod k = 0 \rrbracket \Longrightarrow$$

$$(A \cap [n.., k - \text{Suc } 0]) \odot k = (A \odot k) \cap ([n.., k - \text{Suc } 0] \odot k)$$

<proof>

lemma *iT-Div-assoc:* $I \odot a \odot b = I \odot (a * b)$

<proof>

lemma *iT-Div-commute:* $I \odot a \odot b = I \odot b \odot a$

<proof>

lemma *iT-Mult-Div-self:* $0 < k \Longrightarrow I \otimes k \odot k = I$

<proof>

lemma *iT-Mult-Div:*

$$\llbracket 0 < d; k \bmod d = 0 \rrbracket \Longrightarrow I \otimes k \odot d = I \otimes (k \text{ div } d)$$

<proof>

lemma *iT-Div-Mult-self:*

$$0 < k \Longrightarrow I \odot k \otimes k = \{y. \exists x \in I. y = x - x \bmod k\}$$

<proof>

lemma *iT-Plus-Div-distrib-mod-less*:

$$\forall x \in I. x \bmod m + n \bmod m < m \implies I \oplus n \oslash m = I \oslash m \oplus n \operatorname{div} m$$

<proof>

corollary *iT-Plus-Div-distrib-mod-0*:

$$n \bmod m = 0 \implies I \oplus n \oslash m = I \oslash m \oplus n \operatorname{div} m$$

<proof>

lemma *iT-Div-Min*: $I \neq \{\}$ $\implies iMin (I \oslash k) = iMin I \operatorname{div} k$

<proof>

corollary

iFROM-div-Min: $iMin ([n..] \oslash k) = n \operatorname{div} k$ **and**

iIN-div-Min: $iMin ([n..,d] \oslash k) = n \operatorname{div} k$ **and**

iTILL-div-Min: $iMin ([..n] \oslash k) = 0$ **and**

iMOD-div-Min: $iMin ([r, \bmod m] \oslash k) = r \operatorname{div} k$ **and**

iMODb-div-Min: $iMin ([r, \bmod m, c] \oslash k) = r \operatorname{div} k$

<proof>

lemmas *iT-div-Min* =

iFROM-div-Min

iIN-div-Min

iTILL-div-Min

iMOD-div-Min

iMODb-div-Min

lemma *iT-Div-Max*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies \text{Max } (I \oslash k) = \text{Max } I \operatorname{div} k$

<proof>

corollary

iIN-div-Max: $\text{Max } ([n..,d] \oslash k) = (n + d) \operatorname{div} k$ **and**

iTILL-div-Max: $\text{Max } ([..n] \oslash k) = n \operatorname{div} k$ **and**

iMODb-div-Max: $\text{Max } ([r, \bmod m, c] \oslash k) = (r + m * c) \operatorname{div} k$

<proof>

lemma *iT-Div-0-finite*: $\text{finite } (I \oslash 0)$

<proof>

lemma *iT-Div-infinite-iff*: $0 < k \implies \text{infinite } (I \oslash k) = \text{infinite } I$

<proof>

lemma *iT-Div-finite-iff*: $0 < k \implies \text{finite } (I \oslash k) = \text{finite } I$

<proof>

lemma *iFROM-div*: $0 < k \implies [n..] \oslash k = [n \operatorname{div} k..]$

<proof>

lemma *iIN-div*:

$$0 < k \implies$$

$$[n..,d] \oslash k = [n \operatorname{div} k.., d \operatorname{div} k + (n \bmod k + d \bmod k) \operatorname{div} k]$$

$\langle proof \rangle$

corollary *iIN-div-if*:

$$0 < k \implies [n \dots, d] \odot k = [n \operatorname{div} k \dots, d \operatorname{div} k + (\text{if } n \operatorname{mod} k + d \operatorname{mod} k < k \text{ then } 0 \text{ else } \operatorname{Suc} 0)]$$

$\langle proof \rangle$

corollary *iIN-div-eq1*:

$$\llbracket 0 < k; n \operatorname{mod} k + d \operatorname{mod} k < k \rrbracket \implies [n \dots, d] \odot k = [n \operatorname{div} k \dots, d \operatorname{div} k]$$

$\langle proof \rangle$

corollary *iIN-div-eq2*:

$$\llbracket 0 < k; k \leq n \operatorname{mod} k + d \operatorname{mod} k \rrbracket \implies [n \dots, d] \odot k = [n \operatorname{div} k \dots, \operatorname{Suc} (d \operatorname{div} k)]$$

$\langle proof \rangle$

corollary *iIN-div-mod-eq-0*:

$$\llbracket 0 < k; n \operatorname{mod} k = 0 \rrbracket \implies [n \dots, d] \odot k = [n \operatorname{div} k \dots, d \operatorname{div} k]$$

$\langle proof \rangle$

lemma *iTILL-div*:

$$0 < k \implies [\dots n] \odot k = [\dots n \operatorname{div} k]$$

$\langle proof \rangle$

lemma *iMOD-div-ge*:

$$\llbracket 0 < m; m \leq k \rrbracket \implies [r, \operatorname{mod} m] \odot k = [r \operatorname{div} k \dots]$$

$\langle proof \rangle$

corollary *iMOD-div-self*:

$$0 < m \implies [r, \operatorname{mod} m] \odot m = [r \operatorname{div} m \dots]$$

$\langle proof \rangle$

lemma *iMOD-div*:

$$\llbracket 0 < k; m \operatorname{mod} k = 0 \rrbracket \implies [r, \operatorname{mod} m] \odot k = [r \operatorname{div} k, \operatorname{mod} (m \operatorname{div} k)]$$

$\langle proof \rangle$

lemma *iMODb-div-self*:

$$0 < m \implies [r, \operatorname{mod} m, c] \odot m = [r \operatorname{div} m \dots, c]$$

$\langle proof \rangle$

lemma *iMODb-div-ge*:

$$\llbracket 0 < m; m \leq k \rrbracket \implies [r, \operatorname{mod} m, c] \odot k = [r \operatorname{div} k \dots, (r + m * c) \operatorname{div} k - r \operatorname{div} k]$$

$\langle proof \rangle$

corollary *iMODb-div-ge-if*:

$$\llbracket 0 < m; m \leq k \rrbracket \implies$$

$[r, \text{mod } m, c] \odot k =$
 $[r \text{ div } k \dots, m * c \text{ div } k + (\text{if } r \text{ mod } k + m * c \text{ mod } k < k \text{ then } 0 \text{ else } \text{Suc } 0)]$
 $\langle \text{proof} \rangle$

lemma *iMODb-div*:

$\llbracket 0 < k; m \text{ mod } k = 0 \rrbracket \implies$
 $[r, \text{mod } m, c] \odot k = [r \text{ div } k, \text{mod } (m \text{ div } k), c]$
 $\langle \text{proof} \rangle$

lemmas *iT-div* =

iTILL-div
iFROM-div
iIN-div
iMOD-div
iMODb-div
iT-Div-singleton

This lemma is valid for all $k \leq m$, i. e., also for k with $m \text{ mod } k \neq 0$.

lemma *iMODb-div-unique*:

$\llbracket 0 < k; k \leq m; k \leq c; [r', \text{mod } m', c'] = [r, \text{mod } m, c] \odot k \rrbracket \implies$
 $r' = r \text{ div } k \wedge m' = m \text{ div } k \wedge c' = c$
 $\langle \text{proof} \rangle$

lemma *iMODb-div-mod-gr0-is-0-not-ex0*:

$\llbracket 0 < k; k < m; 0 < m \text{ mod } k; k \leq c; r \text{ mod } k = 0 \rrbracket \implies$
 $\neg(\exists r' m' c'. [r', \text{mod } m', c'] = [r, \text{mod } m, c] \odot k)$
 $\langle \text{proof} \rangle$

lemma *iMODb-div-mod-gr0-not-ex--arith-aux1*:

$\llbracket (0::\text{nat}) < k; k < m; 0 < x1 \rrbracket \implies$
 $x1 * m + x2 - x \text{ mod } k + x3 + x \text{ mod } k = x1 * m + x2 + x3$
 $\langle \text{proof} \rangle$

lemma *iMODb-div-mod-gr0-not-ex*:

$\llbracket 0 < k; k < m; 0 < m \text{ mod } k; k \leq c \rrbracket \implies$
 $\neg(\exists r' m' c'. [r', \text{mod } m', c'] = [r, \text{mod } m, c] \odot k)$
 $\langle \text{proof} \rangle$

lemma *iMOD-div-eq-imp-iMODb-div-eq*:

$\llbracket 0 < k; k \leq m; [r', \text{mod } m'] = [r, \text{mod } m] \odot k \rrbracket \implies$
 $[r', \text{mod } m', c] = [r, \text{mod } m, c] \odot k$
 $\langle \text{proof} \rangle$

lemma *iMOD-div-unique*:

$\llbracket 0 < k; k \leq m; [r', \text{mod } m'] = [r, \text{mod } m] \odot k \rrbracket \implies$
 $r' = r \text{ div } k \wedge m' = m \text{ div } k$

$\langle proof \rangle$

lemma *iMOD-div-mod-gr0-not-ex*:

$\llbracket 0 < k; k < m; 0 < m \bmod k \rrbracket \implies$
 $\neg (\exists r' m'. [r', \bmod m] = [r, \bmod m] \odot k)$
 $\langle proof \rangle$

2.2 Interval cut operators with arithmetic interval operators

lemma

iT-Plus-cut-le2: $(I \oplus k) \downarrow \leq (t + k) = (I \downarrow \leq t) \oplus k$ **and**
iT-Plus-cut-less2: $(I \oplus k) \downarrow < (t + k) = (I \downarrow < t) \oplus k$ **and**
iT-Plus-cut-ge2: $(I \oplus k) \downarrow \geq (t + k) = (I \downarrow \geq t) \oplus k$ **and**
iT-Plus-cut-greater2: $(I \oplus k) \downarrow > (t + k) = (I \downarrow > t) \oplus k$
 $\langle proof \rangle$

lemma *iT-Plus-cut-le*:

$(I \oplus k) \downarrow \leq t = (\text{if } t < k \text{ then } \{\} \text{ else } I \downarrow \leq (t - k) \oplus k)$
 $\langle proof \rangle$

lemma *iT-Plus-cut-less*: $(I \oplus k) \downarrow < t = I \downarrow < (t - k) \oplus k$
 $\langle proof \rangle$

lemma *iT-Plus-cut-ge*: $(I \oplus k) \downarrow \geq t = I \downarrow \geq (t - k) \oplus k$
 $\langle proof \rangle$

lemma *iT-Plus-cut-greater*:

$(I \oplus k) \downarrow > t = (\text{if } t < k \text{ then } I \oplus k \text{ else } I \downarrow > (t - k) \oplus k)$
 $\langle proof \rangle$

lemma

iT-Mult-cut-le2: $0 < k \implies (I \otimes k) \downarrow \leq (t * k) = (I \downarrow \leq t) \otimes k$ **and**
iT-Mult-cut-less2: $0 < k \implies (I \otimes k) \downarrow < (t * k) = (I \downarrow < t) \otimes k$ **and**
iT-Mult-cut-ge2: $0 < k \implies (I \otimes k) \downarrow \geq (t * k) = (I \downarrow \geq t) \otimes k$ **and**
iT-Mult-cut-greater2: $0 < k \implies (I \otimes k) \downarrow > (t * k) = (I \downarrow > t) \otimes k$
 $\langle proof \rangle$

lemma *iT-Mult-cut-le*:

$0 < k \implies (I \otimes k) \downarrow \leq t = (I \downarrow \leq (t \text{ div } k)) \otimes k$
 $\langle proof \rangle$

lemma *iT-Mult-cut-less*:

$0 < k \implies (I \otimes k) \downarrow < t =$
 $(\text{if } t \bmod k = 0 \text{ then } (I \downarrow < (t \text{ div } k)) \text{ else } I \downarrow < \text{Suc } (t \text{ div } k)) \otimes k$
 $\langle proof \rangle$

lemma *iT-Mult-cut-greater*:

$0 < k \implies (I \otimes k) \downarrow > t = (I \downarrow > (t \text{ div } k)) \otimes k$

$\langle proof \rangle$

lemma *iT-Mult-cut-ge*:

$0 < k \implies (I \otimes k) \downarrow_{\geq} t =$
 $(\text{if } t \bmod k = 0 \text{ then } (I \downarrow_{\geq} (t \div k)) \text{ else } I \downarrow_{\geq} \text{Suc } (t \div k)) \otimes k$
 $\langle proof \rangle$

lemma *iT-Plus-neg-cut-le2*: $k \leq t \implies (I \oplus - k) \downarrow_{\leq} (t - k) = (I \downarrow_{\leq} t) \oplus - k$
 $\langle proof \rangle$

lemma *iT-Plus-neg-cut-less2*: $(I \oplus - k) \downarrow_{<} (t - k) = (I \downarrow_{<} t) \oplus - k$
 $\langle proof \rangle$

lemma *iT-Plus-neg-cut-ge2*: $(I \oplus - k) \downarrow_{\geq} (t - k) = (I \downarrow_{\geq} t) \oplus - k$
 $\langle proof \rangle$

lemma *iT-Plus-neg-cut-greater2*: $k \leq t \implies (I \oplus - k) \downarrow_{>} (t - k) = (I \downarrow_{>} t) \oplus - k$
 $\langle proof \rangle$

lemma *iT-Plus-neg-cut-le*: $(I \oplus - k) \downarrow_{\leq} t = I \downarrow_{\leq} (t + k) \oplus - k$
 $\langle proof \rangle$

lemma *iT-Plus-neg-cut-less*: $(I \oplus - k) \downarrow_{<} t = I \downarrow_{<} (t + k) \oplus - k$
 $\langle proof \rangle$

lemma *iT-Plus-neg-cut-ge*: $(I \oplus - k) \downarrow_{\geq} t = I \downarrow_{\geq} (t + k) \oplus - k$
 $\langle proof \rangle$

lemma *iT-Plus-neg-cut-greater*: $(I \oplus - k) \downarrow_{>} t = I \downarrow_{>} (t + k) \oplus - k$
 $\langle proof \rangle$

lemma *iT-Minus-cut-le2*: $t \leq k \implies (k \ominus I) \downarrow_{\leq} (k - t) = k \ominus (I \downarrow_{\geq} t)$
 $\langle proof \rangle$

lemma *iT-Minus-cut-less2*: $(k \ominus I) \downarrow_{<} (k - t) = k \ominus (I \downarrow_{>} t)$
 $\langle proof \rangle$

lemma *iT-Minus-cut-ge2*: $(k \ominus I) \downarrow_{\geq} (k - t) = k \ominus (I \downarrow_{\leq} t)$
 $\langle proof \rangle$

lemma *iT-Minus-cut-greater2*: $t \leq k \implies (k \ominus I) \downarrow_{>} (k - t) = k \ominus (I \downarrow_{<} t)$
 $\langle proof \rangle$

lemma *iT-Minus-cut-le*: $(k \ominus I) \downarrow_{\leq} t = k \ominus (I \downarrow_{\geq} (k - t))$
 $\langle proof \rangle$

lemma *iT-Minus-cut-less*:

$(k \ominus I) \downarrow < t = (\text{if } t \leq k \text{ then } k \ominus (I \downarrow > (k - t)) \text{ else } k \ominus I)$
 $\langle \text{proof} \rangle$

lemma *iT-Minus-cut-ge*:

$(k \ominus I) \downarrow \geq t = (\text{if } t \leq k \text{ then } k \ominus (I \downarrow \leq (k - t)) \text{ else } \{\})$
 $\langle \text{proof} \rangle$

lemma *iT-Minus-cut-greater*: $(k \ominus I) \downarrow > t = k \ominus (I \downarrow < (k - t))$
 $\langle \text{proof} \rangle$

lemma *iT-Div-cut-le*:

$0 < k \implies (I \oslash k) \downarrow \leq t = I \downarrow \leq (t * k + (k - \text{Suc } 0)) \oslash k$
 $\langle \text{proof} \rangle$

lemma *iT-Div-cut-less*:

$0 < k \implies (I \oslash k) \downarrow < t = I \downarrow < (t * k) \oslash k$
 $\langle \text{proof} \rangle$

lemma *iT-Div-cut-ge*:

$0 < k \implies (I \oslash k) \downarrow \geq t = I \downarrow \geq (t * k) \oslash k$
 $\langle \text{proof} \rangle$

lemma *iT-Div-cut-greater*:

$0 < k \implies (I \oslash k) \downarrow > t = I \downarrow > (t * k + (k - \text{Suc } 0)) \oslash k$
 $\langle \text{proof} \rangle$

lemma *iT-Div-cut-le2*:

$0 < k \implies (I \oslash k) \downarrow \leq (t \text{ div } k) = I \downarrow \leq (t - t \text{ mod } k + (k - \text{Suc } 0)) \oslash k$
 $\langle \text{proof} \rangle$

lemma *iT-Div-cut-less2*:

$0 < k \implies (I \oslash k) \downarrow < (t \text{ div } k) = I \downarrow < (t - t \text{ mod } k) \oslash k$
 $\langle \text{proof} \rangle$

lemma *iT-Div-cut-ge2*:

$0 < k \implies (I \oslash k) \downarrow \geq (t \text{ div } k) = I \downarrow \geq (t - t \text{ mod } k) \oslash k$
 $\langle \text{proof} \rangle$

lemma *iT-Div-cut-greater2*:

$0 < k \implies (I \oslash k) \downarrow > (t \text{ div } k) = I \downarrow > (t - t \text{ mod } k + (k - \text{Suc } 0)) \oslash k$
 $\langle \text{proof} \rangle$

2.3 *inext* and *iprev* with interval operators

lemma *iT-Plus-inext*: $\text{inext } (n + k) (I \oplus k) = (\text{inext } n I) + k$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-iprev*: $\text{iprev } (n + k) (I \oplus k) = (\text{iprev } n I) + k$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-inext2*: $k \leq n \implies \text{inext } n (I \oplus k) = (\text{inext } (n - k) I) + k$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-prev2*: $k \leq n \implies \text{iprev } n (I \oplus k) = (\text{iprev } (n - k) I) + k$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-inext*: $\text{inext } (n * k) (I \otimes k) = (\text{inext } n I) * k$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-iprev*: $\text{iprev } (n * k) (I \otimes k) = (\text{iprev } n I) * k$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-inext2-if*:
 $\text{inext } n (I \otimes k) = (\text{if } n \bmod k = 0 \text{ then } (\text{inext } (n \text{ div } k) I) * k \text{ else } n)$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-iprev2-if*:
 $\text{iprev } n (I \otimes k) = (\text{if } n \bmod k = 0 \text{ then } (\text{iprev } (n \text{ div } k) I) * k \text{ else } n)$
 $\langle \text{proof} \rangle$

corollary *iT-Mult-inext2*:
 $n \bmod k = 0 \implies \text{inext } n (I \otimes k) = (\text{inext } (n \text{ div } k) I) * k$
 $\langle \text{proof} \rangle$

corollary *iT-Mult-iprev2*:
 $n \bmod k = 0 \implies \text{iprev } n (I \otimes k) = (\text{iprev } (n \text{ div } k) I) * k$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-inext*:
 $k \leq n \implies \text{inext } (n - k) (I \oplus - k) = \text{inext } n I - k$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-iprev*:
 $\text{iprev } (n - k) (I \oplus - k) = \text{iprev } n (I \downarrow \geq k) - k$
 $\langle \text{proof} \rangle$

corollary *iT-Plus-neg-inext2*: $\text{inext } n (I \oplus - k) = \text{inext } (n + k) I - k$
 $\langle \text{proof} \rangle$

corollary *iT-Plus-neg-iprev2*: $\text{iprev } n (I \oplus - k) = \text{iprev } (n + k) (I \downarrow \geq k) - k$
 $\langle \text{proof} \rangle$

lemma *iT-Minus-inext*:
 $\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{inext } (k - n) (k \ominus I) = k - \text{iprev } n I$
 $\langle \text{proof} \rangle$

corollary *iT-Minus-inext2:*

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{inext } n \ (k \ominus I) = k - \text{iprev } (k - n) \ I$$

<proof>

lemma *iT-Minus-iprev:*

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{iprev } (k - n) \ (k \ominus I) = k - \text{inext } n \ (I \downarrow \leq k)$$

<proof>

lemma *iT-Minus-iprev2:*

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{iprev } n \ (k \ominus I) = k - \text{inext } (k - n) \ (I \downarrow \leq k)$$

<proof>

lemma *iT-Plus-inext-nth:* $I \neq \{\} \implies (I \oplus k) \rightarrow n = (I \rightarrow n) + k$

<proof>

lemma *iT-Plus-iprev-nth:* $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies (I \oplus k) \leftarrow n = (I \leftarrow n) + k$

<proof>

lemma *iT-Mult-inext-nth:* $I \neq \{\} \implies (I \otimes k) \rightarrow n = (I \rightarrow n) * k$

<proof>

lemma *iT-Mult-iprev-nth:* $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies (I \otimes k) \leftarrow n = (I \leftarrow n) * k$

<proof>

lemma *iT-Plus-neg-inext-nth:*

$$I \oplus - \ k \neq \{\} \implies (I \oplus - \ k) \rightarrow n = (I \downarrow \geq k \rightarrow n) - k$$

<proof>

lemma *iT-Plus-neg-iprev-nth:*

$$\llbracket \text{finite } I; I \oplus - \ k \neq \{\} \rrbracket \implies (I \oplus - \ k) \leftarrow n = (I \downarrow \geq k \leftarrow n) - k$$

<proof>

lemma *iT-Minus-inext-nth:*

$$k \ominus I \neq \{\} \implies (k \ominus I) \rightarrow n = k - ((I \downarrow \leq k) \leftarrow n)$$

<proof>

lemma *iT-Minus-iprev-nth:*

$$k \ominus I \neq \{\} \implies (k \ominus I) \leftarrow n = k - ((I \downarrow \leq k) \rightarrow n)$$

<proof>

lemma *iT-Div-ge-inext-nth:*

$$\llbracket I \neq \{\}; \forall x \in I. \forall y \in I. x < y \implies x + k \leq y \rrbracket \implies (I \oslash k) \rightarrow n = (I \rightarrow n) \text{ div } k$$

<proof>

lemma *iT-Div-mod-inext-nth:*

$$\llbracket I \neq \{\}; \forall x \in I. \forall y \in I. x \bmod k = y \bmod k \rrbracket \implies$$

$(I \oslash k) \rightarrow n = (I \rightarrow n) \text{ div } k$
 $\langle \text{proof} \rangle$

lemma *iT-Div-ge-iprev-nth*:
 $\llbracket \text{finite } I; I \neq \{\}; \forall x \in I. \forall y \in I. x < y \longrightarrow x + k \leq y \rrbracket \implies$
 $(I \oslash k) \leftarrow n = (I \leftarrow n) \text{ div } k$
 $\langle \text{proof} \rangle$

lemma *iT-Div-mod-iprev-nth*:
 $\llbracket \text{finite } I; I \neq \{\}; \forall x \in I. \forall y \in I. x \bmod k = y \bmod k \rrbracket \implies$
 $(I \oslash k) \leftarrow n = (I \leftarrow n) \text{ div } k$
 $\langle \text{proof} \rangle$

2.4 Cardinality of intervals with interval operators

lemma *iT-Plus-card*: $\text{card } (I \oplus k) = \text{card } I$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-card*: $0 < k \implies \text{card } (I \otimes k) = \text{card } I$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-card*: $\text{card } (I \oplus - k) = \text{card } (I \downarrow \geq k)$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-card-le*: $\text{card } (I \oplus - k) \leq \text{card } I$
 $\langle \text{proof} \rangle$

lemma *iT-Minus-card*: $\text{card } (k \ominus I) = \text{card } (I \downarrow \leq k)$
 $\langle \text{proof} \rangle$

lemma *iT-Minus-card-le*: $\text{finite } I \implies \text{card } (k \ominus I) \leq \text{card } I$
 $\langle \text{proof} \rangle$

lemma *iT-Div-0-card-if*:
 $\text{card } (I \oslash 0) = (\text{if } I = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$
 $\langle \text{proof} \rangle$

lemma *Int-empty-sum*:
 $(\sum k \leq (n :: \text{nat}). \text{if } \{\} \cap (I \ k) = \{\} \text{ then } 0 \text{ else } \text{Suc } 0) = 0$
 $\langle \text{proof} \rangle$

lemma *iT-Div-mod-partition-card*:
 $\text{card } (I \cap [n * d \dots, d - \text{Suc } 0] \oslash d) =$
 $(\text{if } I \cap [n * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$
 $\langle \text{proof} \rangle$

lemma *iT-Div-conv-count*:
 $0 < d \implies I \oslash d = \{k. I \cap [k * d \dots, d - \text{Suc } 0] \neq \{\}\}$
 $\langle \text{proof} \rangle$

lemma *iT-Div-conv-count2*:

$\llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $I \oslash d = \{k. k \leq n \wedge I \cap [k * d \dots, d - \text{Suc } 0] \neq \{\}\}$
 $\langle \text{proof} \rangle$

lemma *mod-partition-count-Suc*:

$\{k. k \leq \text{Suc } n \wedge I \cap [k * d \dots, d - \text{Suc } 0] \neq \{\}\} =$
 $\{k. k \leq n \wedge I \cap [k * d \dots, d - \text{Suc } 0] \neq \{\}\} \cup$
 $(\text{if } I \cap [\text{Suc } n * d \dots, d - \text{Suc } 0] \neq \{\} \text{ then } \{\text{Suc } n\} \text{ else } \{\})$
 $\langle \text{proof} \rangle$

lemma *iT-Div-card*:

$\bigwedge I. \llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $\text{card } (I \oslash d) = (\sum k \leq n.$
 $\text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$
 $\langle \text{proof} \rangle$

corollary *iT-Div-card-Suc*:

$\bigwedge I. \llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $\text{card } (I \oslash d) = (\sum k < \text{Suc } n.$
 $\text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$
 $\langle \text{proof} \rangle$

corollary *iT-Div-Max-card*: $\llbracket 0 < d; \text{finite } I \rrbracket \implies$

$\text{card } (I \oslash d) = (\sum k \leq \text{Max } I \text{ div } d.$
 $\text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$
 $\langle \text{proof} \rangle$

lemma *iT-Div-card-le*: $0 < k \implies \text{card } (I \oslash k) \leq \text{card } I$

$\langle \text{proof} \rangle$

lemma *iT-Div-card-inj-on*:

$\text{inj-on } (\lambda n. n \text{ div } k) I \implies \text{card } (I \oslash k) = \text{card } I$
 $\langle \text{proof} \rangle$

lemma *mod-Suc'*:

$0 < n \implies \text{Suc } m \text{ mod } n = (\text{if } m \text{ mod } n < n - \text{Suc } 0 \text{ then } \text{Suc } (m \text{ mod } n) \text{ else } 0)$
 $\langle \text{proof} \rangle$

lemma *div-Suc*:

$0 < n \implies \text{Suc } m \text{ div } n = (\text{if } \text{Suc } (m \text{ mod } n) = n \text{ then } \text{Suc } (m \text{ div } n) \text{ else } m \text{ div } n)$
 $\langle \text{proof} \rangle$

lemma *div-Suc'*:

$0 < n \implies \text{Suc } m \text{ div } n = (\text{if } m \bmod n < n - \text{Suc } 0 \text{ then } m \text{ div } n \text{ else } \text{Suc } (m \text{ div } n))$
 $\langle \text{proof} \rangle$

lemma *iT-Div-card-ge-aux*:

$\bigwedge I. \llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $\text{card } I \text{ div } d + (\text{if } \text{card } I \bmod d = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq \text{card } (I \oslash d)$
 $\langle \text{proof} \rangle$

lemma *iT-Div-card-ge*:

$\text{card } I \text{ div } d + (\text{if } \text{card } I \bmod d = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq \text{card } (I \oslash d)$
 $\langle \text{proof} \rangle$

corollary *iT-Div-card-ge-div*: $\text{card } I \text{ div } d \leq \text{card } (I \oslash d)$

$\langle \text{proof} \rangle$

There is no better lower bound function f for $i \oslash d$ with $\text{card } i$ and d as arguments.

lemma *iT-Div-card-ge--is-maximal-lower-bound*:

$\forall I d. \text{card } I \text{ div } d + (\text{if } \text{card } I \bmod d = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq f (\text{card } I) d \wedge$
 $f (\text{card } I) d \leq \text{card } (I \oslash d) \implies$
 $f (\text{card } (I :: \text{nat set})) d = \text{card } I \text{ div } d + (\text{if } \text{card } I \bmod d = 0 \text{ then } 0 \text{ else } \text{Suc } 0)$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-icard*: $\text{icard } (I \oplus k) = \text{icard } I$

$\langle \text{proof} \rangle$

lemma *iT-Mult-icard*: $0 < k \implies \text{icard } (I \otimes k) = \text{icard } I$

$\langle \text{proof} \rangle$

lemma *iT-Plus-neg-icard*: $\text{icard } (I \oplus - k) = \text{icard } (I \downarrow \geq k)$

$\langle \text{proof} \rangle$

lemma *iT-Plus-neg-icard-le*: $\text{icard } (I \oplus - k) \leq \text{icard } I$

$\langle \text{proof} \rangle$

lemma *iT-Minus-icard*: $\text{icard } (k \ominus I) = \text{icard } (I \downarrow \leq k)$

$\langle \text{proof} \rangle$

lemma *iT-Minus-icard-le*: $\text{icard } (k \ominus I) \leq \text{icard } I$

$\langle \text{proof} \rangle$

lemma *iT-Div-0-icard-if*: $\text{icard } (I \oslash 0) = \text{enat } (\text{if } I = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$

$\langle \text{proof} \rangle$

lemma *iT-Div-mod-partition-icard*:

$\text{icard } (I \cap [n * d .. d - \text{Suc } 0] \oslash d) =$

enat (if $I \cap [n * d \dots, d - \text{Suc } 0] = \{\}$ then 0 else $\text{Suc } 0$)
 <proof>

lemma *iT-Div-icard*:

$\llbracket 0 < d; \text{finite } I \implies \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $\text{icard } (I \oslash d) =$
 (if *finite* I then *enat* ($\sum_{k \leq n} \text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\}$ then 0 else $\text{Suc } 0$) else ∞)
 <proof>

corollary *iT-Div-Max-icard*: $0 < d \implies$

$\text{icard } (I \oslash d) = (\text{if } \text{finite } I$
 then *enat* ($\sum_{k \leq \text{Max } I \text{ div } d} \text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\}$ then 0 else $\text{Suc } 0$) else ∞)
 <proof>

lemma *iT-Div-icard-le*: $0 < k \implies \text{icard } (I \oslash k) \leq \text{icard } I$
 <proof>

lemma *iT-Div-icard-inj-on*: $\text{inj-on } (\lambda n. n \text{ div } k) I \implies \text{icard } (I \oslash k) = \text{icard } I$
 <proof>

lemma *iT-Div-icard-ge*: $\text{icard } I \text{ div } (\text{enat } d) + \text{enat } (\text{if } \text{icard } I \bmod (\text{enat } d) = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq \text{icard } (I \oslash d)$
 <proof>

corollary *iT-Div-icard-ge-div*: $\text{icard } I \text{ div } (\text{enat } d) \leq \text{icard } (I \oslash d)$
 <proof>

lemma *iT-Div-icard-ge-is-maximal-lower-bound*:

$\forall I d. \text{icard } I \text{ div } (\text{enat } d) + \text{enat } (\text{if } \text{icard } I \bmod (\text{enat } d) = 0 \text{ then } 0 \text{ else } \text{Suc } 0)$
 $\leq f (\text{icard } I) d \wedge$
 $f (\text{icard } I) d \leq \text{icard } (I \oslash d) \implies$
 $f (\text{icard } (I :: \text{nat set})) d =$
 $\text{icard } I \text{ div } (\text{enat } d) + \text{enat } (\text{if } \text{icard } I \bmod (\text{enat } d) = 0 \text{ then } 0 \text{ else } \text{Suc } 0)$
 <proof>

2.5 Results about sets of intervals

2.5.1 Set of intervals without and with empty interval

definition *iFROM-UN-set* :: (nat set) set
 where *iFROM-UN-set* $\equiv \bigcup n. \{[n \dots]\}$

definition *iTILL-UN-set* :: (nat set) set
 where *iTILL-UN-set* $\equiv \bigcup n. \{[\dots n]\}$

definition *iIN-UN-set* :: (nat set) set
 where *iIN-UN-set* $\equiv \bigcup n d. \{[n \dots, d]\}$

definition *iMOD-UN-set* :: (nat set) set
 where *iMOD-UN-set* $\equiv \bigcup r\ m. \{[r, \text{mod } m]\}$

definition *iMODb-UN-set* :: (nat set) set
 where *iMODb-UN-set* $\equiv \bigcup r\ m\ c. \{[r, \text{mod } m, c]\}$

definition *iFROM-set* :: (nat set) set
 where *iFROM-set* $\equiv \{[n..] \mid n. \text{True}\}$

definition *iTILL-set* :: (nat set) set
 where *iTILL-set* $\equiv \{[..n] \mid n. \text{True}\}$

definition *iIN-set* :: (nat set) set
 where *iIN-set* $\equiv \{[n..,d] \mid n\ d. \text{True}\}$

definition *iMOD-set* :: (nat set) set
 where *iMOD-set* $\equiv \{[r, \text{mod } m] \mid r\ m. \text{True}\}$

definition *iMODb-set* :: (nat set) set
 where *iMODb-set* $\equiv \{[r, \text{mod } m, c] \mid r\ m\ c. \text{True}\}$

definition *iMOD2-set* :: (nat set) set
 where *iMOD2-set* $\equiv \{[r, \text{mod } m] \mid r\ m. 2 \leq m\}$

definition *iMODb2-set* :: (nat set) set
 where *iMODb2-set* $\equiv \{[r, \text{mod } m, c] \mid r\ m\ c. 2 \leq m \wedge 1 \leq c\}$

definition *iMOD2-UN-set* :: (nat set) set
 where *iMOD2-UN-set* $\equiv \bigcup r. \bigcup m \in \{2..\}. \{[r, \text{mod } m]\}$

definition *iMODb2-UN-set* :: (nat set) set
 where *iMODb2-UN-set* $\equiv \bigcup r. \bigcup m \in \{2..\}. \bigcup c \in \{1..\}. \{[r, \text{mod } m, c]\}$

lemmas *i-set-defs* =
iFROM-set-def *iTILL-set-def* *iIN-set-def*
iMOD-set-def *iMODb-set-def*
iMOD2-set-def *iMODb2-set-def*

lemmas *i-UN-set-defs* =
iFROM-UN-set-def *iTILL-UN-set-def* *iIN-UN-set-def*
iMOD-UN-set-def *iMODb-UN-set-def*
iMOD2-UN-set-def *iMODb2-UN-set-def*

lemma *iFROM-set-UN-set-eq*: *iFROM-set* = *iFROM-UN-set*
 <proof>

lemma

iTILL-set-UN-set-eq: iTILL-set = iTILL-UN-set and
iIN-set-UN-set-eq: iIN-set = iIN-UN-set and
iMOD-set-UN-set-eq: iMOD-set = iMOD-UN-set and
iMODb-set-UN-set-eq: iMODb-set = iMODb-UN-set

<proof>

lemma *iMOD2-set-UN-set-eq: iMOD2-set = iMOD2-UN-set*

<proof>

lemma *iMODb2-set-UN-set-eq: iMODb2-set = iMODb2-UN-set*

<proof>

lemmas *i-set-i-UN-set-sets-eq =*

iFROM-set-UN-set-eq
iTILL-set-UN-set-eq
iIN-set-UN-set-eq
iMOD-set-UN-set-eq
iMODb-set-UN-set-eq
iMOD2-set-UN-set-eq
iMODb2-set-UN-set-eq

lemma *iMOD2-set-iMOD-set-subset: iMOD2-set $\subseteq$ iMOD-set*

<proof>

lemma *iMODb2-set-iMODb-set-subset: iMODb2-set $\subseteq$ iMODb-set*

<proof>

definition *i-set :: (nat set) set*

where *i-set $\equiv$ iFROM-set $\cup$ iTILL-set $\cup$ iIN-set $\cup$ iMOD-set $\cup$ iMODb-set*

definition *i-UN-set :: (nat set) set*

where *i-UN-set $\equiv$ iFROM-UN-set $\cup$ iTILL-UN-set $\cup$ iIN-UN-set $\cup$ iMOD-UN-set $\cup$ iMODb-UN-set*

Minimal definitions for *i-set* and *i-set*

definition *i-set-min :: (nat set) set*

where *i-set-min $\equiv$ iFROM-set $\cup$ iIN-set $\cup$ iMOD2-set $\cup$ iMODb2-set*

definition *i-UN-set-min :: (nat set) set*

where *i-UN-set-min $\equiv$ iFROM-UN-set $\cup$ iIN-UN-set $\cup$ iMOD2-UN-set $\cup$ iMODb2-UN-set*

definition *i-set0 :: (nat set) set*

where *i-set0 $\equiv$ insert {} i-set*

lemma *i-set-i-UN-set-eq: i-set = i-UN-set*

<proof>

lemma *i-set-min-i-UN-set-min-eq*: $i\text{-set-min} = i\text{-UN-set-min}$
 <proof>

lemma *i-set-min-eq*: $i\text{-set} = i\text{-set-min}$
 <proof>

corollary *i-UN-set-i-UN-min-set-eq*: $i\text{-UN-set} = i\text{-UN-set-min}$
 <proof>

lemma *i-set-min-is-minimal-let*:
 let $s1 = i\text{FROM-set}$; $s2 = i\text{IN-set}$; $s3 = i\text{MOD2-set}$; $s4 = i\text{MODb2-set}$ in
 $s1 \cap s2 = \{\} \wedge s1 \cap s3 = \{\} \wedge s1 \cap s4 = \{\} \wedge$
 $s2 \cap s3 = \{\} \wedge s2 \cap s4 = \{\} \wedge s3 \cap s4 = \{\}$
 <proof>

lemmas *i-set-min-is-minimal* = *i-set-min-is-minimal-let*[simplified]

inductive-set *i-set-ind*:: (nat set) set
 where

i-set-ind-FROM[intro!]: $[n \dots] \in i\text{-set-ind}$
i-set-ind-TILL[intro!]: $[\dots n] \in i\text{-set-ind}$
i-set-ind-IN[intro!]: $[n \dots, d] \in i\text{-set-ind}$
i-set-ind-MOD[intro!]: $[r, \text{mod } m] \in i\text{-set-ind}$
i-set-ind-MODb[intro!]: $[r, \text{mod } m, c] \in i\text{-set-ind}$

inductive-set *i-set0-ind*:: (nat set) set
 where

i-set0-ind-empty[intro!]: $\{\} \in i\text{-set0-ind}$
i-set0-ind-i-set[intro]: $I \in i\text{-set-ind} \implies I \in i\text{-set0-ind}$

The introduction rule *i-set0-ind-i-set* is not declared a safe introduction rule, because it would disturb the correct usage of the *safe* method.

lemma *i-set-ind-subset-i-set0-ind*: $i\text{-set-ind} \subseteq i\text{-set0-ind}$
 <proof>

lemma
i-set0-ind-FROM[intro!]: $[n \dots] \in i\text{-set0-ind}$ **and**
i-set0-ind-TILL[intro!]: $[\dots n] \in i\text{-set0-ind}$ **and**
i-set0-ind-IN[intro!]: $[n \dots, d] \in i\text{-set0-ind}$ **and**
i-set0-ind-MOD[intro!]: $[r, \text{mod } m] \in i\text{-set0-ind}$ **and**
i-set0-ind-MODb[intro!]: $[r, \text{mod } m, c] \in i\text{-set0-ind}$
 <proof>

lemmas *i-set0-ind-intros2* =
i-set0-ind-empty
i-set0-ind-FROM
i-set0-ind-TILL

i-set0-ind-IN
i-set0-ind-MOD
i-set0-ind-MODb

lemma *i-set-i-set-ind-eq*: $i\text{-set} = i\text{-set-ind}$
 $\langle \text{proof} \rangle$

lemma *i-set0-i-set0-ind-eq*: $i\text{-set0} = i\text{-set0-ind}$
 $\langle \text{proof} \rangle$

lemma *i-set-imp-not-empty*: $I \in i\text{-set} \implies I \neq \{\}$
 $\langle \text{proof} \rangle$

lemma *i-set0-i-set-mem-conv*: $(I \in i\text{-set0}) = (I \in i\text{-set} \vee I = \{\})$
 $\langle \text{proof} \rangle$

lemma *i-set-i-set0-mem-conv*: $(I \in i\text{-set}) = (I \in i\text{-set0} \wedge I \neq \{\})$
 $\langle \text{proof} \rangle$

lemma *i-set0-i-set-conv*: $i\text{-set0} - \{\{\}\} = i\text{-set}$
 $\langle \text{proof} \rangle$

corollary *i-set-subset-i-set0*: $i\text{-set} \subseteq i\text{-set0}$
 $\langle \text{proof} \rangle$

lemma *i-set-singleton*: $\{a\} \in i\text{-set}$
 $\langle \text{proof} \rangle$

lemma *i-set0-singleton*: $\{a\} \in i\text{-set0}$
 $\langle \text{proof} \rangle$

corollary
i-set-FROM[intro!]: $[n \dots] \in i\text{-set}$ **and**
i-set-TILL[intro!]: $[\dots n] \in i\text{-set}$ **and**
i-set-IN[intro!]: $[n \dots, d] \in i\text{-set}$ **and**
i-set-MOD[intro!]: $[r, \text{mod } m] \in i\text{-set}$ **and**
i-set-MODb[intro!]: $[r, \text{mod } m, c] \in i\text{-set}$
 $\langle \text{proof} \rangle$

lemmas *i-set-intros* =
i-set-FROM
i-set-TILL
i-set-IN
i-set-MOD
i-set-MODb

lemma
i-set0-empty[intro!]: $\{\} \in i\text{-set0}$ **and**
i-set0-FROM[intro!]: $[n \dots] \in i\text{-set0}$ **and**

$i\text{-set0-TILL}[\text{intro!}] : [\dots n] \in i\text{-set0} \text{ and}$
 $i\text{-set0-IN}[\text{intro!}] : [n \dots, d] \in i\text{-set0} \text{ and}$
 $i\text{-set0-MOD}[\text{intro!}] : [r, \text{mod } m] \in i\text{-set0} \text{ and}$
 $i\text{-set0-MODb}[\text{intro!}] : [r, \text{mod } m, c] \in i\text{-set0}$
 $\langle \text{proof} \rangle$

lemmas $i\text{-set0-intros} =$

$i\text{-set0-empty}$
 $i\text{-set0-FROM}$
 $i\text{-set0-TILL}$
 $i\text{-set0-IN}$
 $i\text{-set0-MOD}$
 $i\text{-set0-MODb}$

lemma $i\text{-set-infinite-as-iMOD}$:

$\llbracket I \in i\text{-set}; \text{infinite } I \rrbracket \implies \exists r \ m. I = [r, \text{mod } m]$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-finite-as-iMODb}$:

$\llbracket I \in i\text{-set}; \text{finite } I \rrbracket \implies \exists r \ m \ c. I = [r, \text{mod } m, c]$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-as-iMOD-iMODb}$:

$I \in i\text{-set} \implies (\exists r \ m. I = [r, \text{mod } m]) \vee (\exists r \ m \ c. I = [r, \text{mod } m, c])$
 $\langle \text{proof} \rangle$

2.5.2 Interval sets are closed under cutting

lemma $i\text{-set-cut-le-ge-closed-disj}$:

$\llbracket I \in i\text{-set}; t \in I; \text{cut-op} = (\downarrow \leq) \vee \text{cut-op} = (\downarrow \geq) \rrbracket \implies$
 $\text{cut-op } I \ t \in i\text{-set}$
 $\langle \text{proof} \rangle$

corollary

$i\text{-set-cut-le-closed}: \llbracket I \in i\text{-set}; t \in I \rrbracket \implies I \downarrow \leq t \in i\text{-set} \text{ and}$
 $i\text{-set-cut-ge-closed}: \llbracket I \in i\text{-set}; t \in I \rrbracket \implies I \downarrow \geq t \in i\text{-set}$
 $\langle \text{proof} \rangle$

lemmas $i\text{-set-cut-le-ge-closed} = i\text{-set-cut-le-closed } i\text{-set-cut-ge-closed}$

lemma $i\text{-set0-cut-closed-disj}$:

$\llbracket I \in i\text{-set0};$
 $\text{cut-op} = (\downarrow \leq) \vee \text{cut-op} = (\downarrow \geq) \vee$
 $\text{cut-op} = (\downarrow <) \vee \text{cut-op} = (\downarrow >) \rrbracket \implies$
 $\text{cut-op } I \ t \in i\text{-set0}$
 $\langle \text{proof} \rangle$

corollary

i-set0-cut-le-closed: $I \in i\text{-set0} \implies I \downarrow \leq t \in i\text{-set0}$ and
i-set0-cut-less-closed: $I \in i\text{-set0} \implies I \downarrow < t \in i\text{-set0}$ and
i-set0-cut-ge-closed: $I \in i\text{-set0} \implies I \downarrow \geq t \in i\text{-set0}$ and
i-set0-cut-greater-closed: $I \in i\text{-set0} \implies I \downarrow > t \in i\text{-set0}$
 <proof>

lemmas *i-set0-cut-closed* =
i-set0-cut-le-closed
i-set0-cut-less-closed
i-set0-cut-ge-closed
i-set0-cut-greater-closed

2.5.3 Interval sets are closed under addition and multiplication

lemma *i-set-Plus-closed*: $I \in i\text{-set} \implies I \oplus k \in i\text{-set}$
 <proof>

lemma *i-set-Mult-closed*: $I \in i\text{-set} \implies I \otimes k \in i\text{-set}$
 <proof>

lemma *i-set0-Plus-closed*: $I \in i\text{-set0} \implies I \oplus k \in i\text{-set0}$
 <proof>

lemma *i-set0-Mult-closed*: $I \in i\text{-set0} \implies I \otimes k \in i\text{-set0}$
 <proof>

2.5.4 Interval sets are closed with certain conditions under subtraction

lemma *i-set-Plus-neg-closed*:
 $\llbracket I \in i\text{-set}; \exists x \in I. k \leq x \rrbracket \implies I \oplus - k \in i\text{-set}$
 <proof>

lemma *i-set-Minus-closed*:
 $\llbracket I \in i\text{-set}; i\text{Min } I \leq k \rrbracket \implies k \ominus I \in i\text{-set}$
 <proof>

lemma *i-set0-Plus-neg-closed*: $I \in i\text{-set0} \implies I \oplus - k \in i\text{-set0}$
 <proof>

lemma *i-set0-Minus-closed*: $I \in i\text{-set0} \implies k \ominus I \in i\text{-set0}$
 <proof>

lemmas *i-set-IntOp-closed* =
i-set-Plus-closed
i-set-Mult-closed
i-set-Plus-neg-closed
i-set-Minus-closed

lemmas *i-set0-IntOp-closed* =
i-set0-Plus-closed
i-set0-Mult-closed
i-set0-Plus-neg-closed
i-set0-Minus-closed

2.5.5 Interval sets are not closed under division

lemma *iMOD-div-mod-gr0-not-in-i-set*:
 $\llbracket 0 < k; k < m; 0 < m \bmod k \rrbracket \implies [r, \bmod m] \odot k \notin i\text{-set}$
 $\langle \text{proof} \rangle$

lemma *iMODb-div-mod-gr0-not-in-i-set*:
 $\llbracket 0 < k; k < m; 0 < m \bmod k; k \leq c \rrbracket \implies [r, \bmod m, c] \odot k \notin i\text{-set}$
 $\langle \text{proof} \rangle$

lemma $[0, \bmod 3] \odot 2 \notin i\text{-set}$
 $\langle \text{proof} \rangle$

lemma *i-set-Div-not-closed*: $\text{Suc } 0 < k \implies \exists I \in i\text{-set}. I \odot k \notin i\text{-set}$
 $\langle \text{proof} \rangle$

lemma *i-set0-Div-not-closed*: $\text{Suc } 0 < k \implies \exists I \in i\text{-set0}. I \odot k \notin i\text{-set0}$
 $\langle \text{proof} \rangle$

2.5.6 Sets of intervals closed under division

inductive-set *NatMultiples* :: *nat set* $\Rightarrow$ *nat set*
for *F* :: *nat set*

where

NatMultiples-Factor: $k \in F \implies k \in \text{NatMultiples } F$
NatMultiples-Product: $\llbracket k \in F; m \in \text{NatMultiples } F \rrbracket \implies k * m \in \text{NatMultiples } F$

lemma *NatMultiples-ex-divisor*: $m \in \text{NatMultiples } F \implies \exists k \in F. m \bmod k = 0$
 $\langle \text{proof} \rangle$

lemma *NatMultiples-product-factor*: $\llbracket a \in F; b \in F \rrbracket \implies a * b \in \text{NatMultiples } F$
 $\langle \text{proof} \rangle$

lemma *NatMultiples-product-factor-multiple*:
 $\llbracket a \in F; b \in \text{NatMultiples } F \rrbracket \implies a * b \in \text{NatMultiples } F$
 $\langle \text{proof} \rangle$

lemma *NatMultiples-product-multiple-factor*:
 $\llbracket a \in \text{NatMultiples } F; b \in F \rrbracket \implies a * b \in \text{NatMultiples } F$
 $\langle \text{proof} \rangle$

lemma *NatMultiples-product-multiple*:
 $\llbracket a \in \text{NatMultiples } F; b \in \text{NatMultiples } F \rrbracket \implies a * b \in \text{NatMultiples } F$
 $\langle \text{proof} \rangle$

Subset of *i-set* containing only continuous intervals, i. e., without *iMOD* and *iMODb*.

inductive-set *i-set-cont* :: (nat set) set
where

i-set-cont-FROM[intro]: $[n..] \in i\text{-set-cont}$
 | *i-set-cont-TILL*[intro]: $[\dots n] \in i\text{-set-cont}$
 | *i-set-cont-IN*[intro]: $[n.., d] \in i\text{-set-cont}$

definition *i-set0-cont* :: (nat set) set
where *i-set0-cont* $\equiv$ insert {} *i-set-cont*

lemma *i-set-cont-subset-i-set*: *i-set-cont* $\subseteq$ *i-set*
 <proof>

lemma *i-set0-cont-subset-i-set0*: *i-set0-cont* $\subseteq$ *i-set0*
 <proof>

Minimal definition of *i-set-cont*

inductive-set *i-set-cont-min* :: (nat set) set
where

i-set-cont-FROM[intro]: $[n..] \in i\text{-set-cont-min}$
 | *i-set-cont-IN*[intro]: $[n.., d] \in i\text{-set-cont-min}$

definition *i-set0-cont-min* :: (nat set) set
where *i-set0-cont-min* $\equiv$ insert {} *i-set-cont-min*

lemma *i-set-cont-min-eq*: *i-set-cont* = *i-set-cont-min*
 <proof>

inext and *iprev* with continuous intervals

lemma *i-set-cont-inext*:
 $\llbracket I \in i\text{-set-cont}; n \in I; \text{finite } I \implies n < \text{Max } I \rrbracket \implies \text{inext } n \ I = \text{Suc } n$
 <proof>

lemma *i-set-cont-iprev*:
 $\llbracket I \in i\text{-set-cont}; n \in I; i\text{Min } I < n \rrbracket \implies \text{iprev } n \ I = n - \text{Suc } 0$
 <proof>

lemma *i-set-cont-inext--less*:
 $\llbracket I \in i\text{-set-cont}; n \in I; n < n0; n0 \in I \rrbracket \implies \text{inext } n \ I = \text{Suc } n$
 <proof>

lemma *i-set-cont-iprev--greater*:
 $\llbracket I \in i\text{-set-cont}; n \in I; n0 < n; n0 \in I \rrbracket \implies \text{iprev } n \ I = n - \text{Suc } 0$
 <proof>

Sets of modulo intervals

inductive-set *i-set-mult* :: nat $\Rightarrow$ (nat set) set

for $k :: \text{nat}$
where
 $i\text{-set-mult-FROM}[intro!]: [n..] \in i\text{-set-mult } k$
 $| i\text{-set-mult-TILL}[intro!]: [\dots n] \in i\text{-set-mult } k$
 $| i\text{-set-mult-IN}[intro!]: [n.., d] \in i\text{-set-mult } k$
 $| i\text{-set-mult-MOD}[intro!]: [r, \text{mod } m * k] \in i\text{-set-mult } k$
 $| i\text{-set-mult-MODb}[intro!]: [r, \text{mod } m * k, c] \in i\text{-set-mult } k$

definition $i\text{-set0-mult} :: \text{nat} \Rightarrow (\text{nat set}) \text{ set}$
where $i\text{-set0-mult } k \equiv \text{insert } \{\} (i\text{-set-mult } k)$

lemma
 $i\text{-set0-mult-empty}[intro!]: \{\} \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-FROM}[intro!]: [n..] \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-TILL}[intro!]: [\dots n] \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-IN}[intro!]: [n.., d] \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-MOD}[intro!]: [r, \text{mod } m * k] \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-MODb}[intro!]: [r, \text{mod } m * k, c] \in i\text{-set0-mult } k$
 $\langle \text{proof} \rangle$

lemmas $i\text{-set0-mult-intros} =$
 $i\text{-set0-mult-empty}$
 $i\text{-set0-mult-FROM}$
 $i\text{-set0-mult-TILL}$
 $i\text{-set0-mult-IN}$
 $i\text{-set0-mult-MOD}$
 $i\text{-set0-mult-MODb}$

lemma $i\text{-set-mult-subset-i-set0-mult}: i\text{-set-mult } k \subseteq i\text{-set0-mult } k$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-mult-subset-i-set}: i\text{-set-mult } k \subseteq i\text{-set}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set0-mult-subset-i-set0}: i\text{-set0-mult } k \subseteq i\text{-set0}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-mult-0-eq-i-set-cont}: i\text{-set-mult } 0 = i\text{-set-cont}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set0-mult-0-eq-i-set0-cont}: i\text{-set0-mult } 0 = i\text{-set0-cont}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-mult-1-eq-i-set}: i\text{-set-mult } (\text{Suc } 0) = i\text{-set}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set0-mult-1-eq-i-set0}: i\text{-set0-mult } (\text{Suc } 0) = i\text{-set0}$
 $\langle \text{proof} \rangle$

lemma *i-set-mult-imp-not-empty*: $I \in i\text{-set-mult } k \implies I \neq \{\}$
 $\langle \text{proof} \rangle$

lemma *iMOD-in-i-set-mult-imp-divisor-mod-0*:
 $\llbracket m \neq \text{Suc } 0; [r, \text{mod } m] \in i\text{-set-mult } k \rrbracket \implies m \bmod k = 0$
 $\langle \text{proof} \rangle$

lemma
divisor-mod-0-imp-iMOD-in-i-set-mult: $m \bmod k = 0 \implies [r, \text{mod } m] \in i\text{-set-mult } k$ **and**
divisor-mod-0-imp-iMODb-in-i-set-mult: $m \bmod k = 0 \implies [r, \text{mod } m, c] \in i\text{-set-mult } k$
 $\langle \text{proof} \rangle$

lemma *iMOD-in-i-set-mult--divisor-mod-0-conv*:
 $m \neq \text{Suc } 0 \implies ([r, \text{mod } m] \in i\text{-set-mult } k) = (m \bmod k = 0)$
 $\langle \text{proof} \rangle$

lemma *i-set-mult-neq1-subset-i-set*: $k \neq \text{Suc } 0 \implies i\text{-set-mult } k \subset i\text{-set}$
 $\langle \text{proof} \rangle$

lemma *mod-0-imp-i-set-mult-subset*:
 $a \bmod b = 0 \implies i\text{-set-mult } a \subseteq i\text{-set-mult } b$
 $\langle \text{proof} \rangle$

lemma *i-set-mult-subset-imp-mod-0*:
 $\llbracket a \neq \text{Suc } 0; (i\text{-set-mult } a \subseteq i\text{-set-mult } b) \rrbracket \implies a \bmod b = 0$
 $\langle \text{proof} \rangle$

lemma *i-set-mult-subset-conv*:
 $a \neq \text{Suc } 0 \implies (i\text{-set-mult } a \subseteq i\text{-set-mult } b) = (a \bmod b = 0)$
 $\langle \text{proof} \rangle$

lemma *i-set-mult-mod-0-div*:
 $\llbracket I \in i\text{-set-mult } k; k \bmod d = 0 \rrbracket \implies I \oslash d \in i\text{-set-mult } (k \text{ div } d)$
 $\langle \text{proof} \rangle$

Intervals from *i-set-mult* k remain in *i-set* after division by d a divisor of k .

corollary *i-set-mult-mod-0-div-i-set*:
 $\llbracket I \in i\text{-set-mult } k; k \bmod d = 0 \rrbracket \implies I \oslash d \in i\text{-set}$
 $\langle \text{proof} \rangle$

corollary *i-set-mult-div-self-i-set*:
 $I \in i\text{-set-mult } k \implies I \oslash k \in i\text{-set}$
 $\langle \text{proof} \rangle$

lemma *i-set-mod-0-mult-in-i-set-mult*:
 $\llbracket I \in i\text{-set}; m \bmod k = 0 \rrbracket \implies I \otimes m \in i\text{-set-mult } k$

$\langle proof \rangle$

lemma *i-set-self-mult-in-i-set-mult*:

$I \in i\text{-set} \implies I \otimes k \in i\text{-set-mult } k$

$\langle proof \rangle$

lemma *i-set-mult-mod-gr0-div-not-in-i-set*:

$\llbracket 0 < k; 0 < d; 0 < k \bmod d \rrbracket \implies \exists I \in i\text{-set-mult } k. I \oslash d \notin i\text{-set}$

$\langle proof \rangle$

lemma *i-set0-mult-mod-0-div*:

$\llbracket I \in i\text{-set0-mult } k; k \bmod d = 0 \rrbracket \implies I \oslash d \in i\text{-set0-mult } (k \text{ div } d)$

$\langle proof \rangle$

corollary *i-set0-mult-mod-0-div-i-set0*:

$\llbracket I \in i\text{-set0-mult } k; k \bmod d = 0 \rrbracket \implies I \oslash d \in i\text{-set0}$

$\langle proof \rangle$

corollary *i-set0-mult-div-self-i-set0*:

$I \in i\text{-set0-mult } k \implies I \oslash k \in i\text{-set0}$

$\langle proof \rangle$

lemma *i-set0-mod-0-mult-in-i-set0-mult*:

$\llbracket I \in i\text{-set0}; m \bmod k = 0 \rrbracket \implies I \otimes m \in i\text{-set0-mult } k$

$\langle proof \rangle$

lemma *i-set0-self-mult-in-i-set0-mult*:

$I \in i\text{-set0} \implies I \otimes k \in i\text{-set0-mult } k$

$\langle proof \rangle$

lemma *i-set0-mult-mod-gr0-div-not-in-i-set0*:

$\llbracket 0 < k; 0 < d; 0 < k \bmod d \rrbracket \implies \exists I \in i\text{-set0-mult } k. I \oslash d \notin i\text{-set0}$

$\langle proof \rangle$

end

3 Temporal logic operators on natural intervals

theory *IL-TemporalOperators*

imports *IL-IntervalOperators*

begin

Bool : some additional properties

instantiation *bool* :: {ord, zero, one, plus, times, order}

begin

definition *Zero-bool-def* [simp]: $0 \equiv \text{False}$

definition *One-bool-def* [simp]: $1 \equiv \text{True}$

definition *add-bool-def*: $a + b \equiv a \vee b$

definition *mult-bool-def*: $a * b \equiv a \wedge b$

instance $\langle proof \rangle$

end

value *False* < *True*

value *True* < *True*

value *True* ≤ *True*

lemmas *bool-op-rel-defs* =

add-bool-def

mult-bool-def

less-bool-def

le-bool-def

instance *bool* :: *wellorder*

$\langle proof \rangle$

instance *bool* :: *comm-semiring-1*

$\langle proof \rangle$

3.1 Basic definitions

lemma *UNIV-nat*: $\mathbb{N} = (UNIV::nat\ set)$

$\langle proof \rangle$

Universal temporal operator: Always/Globally

definition *iAll* :: $iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Always

where *iAll* *I* *P* $\equiv \forall t \in I. P\ t$

Existential temporal operator: Eventually/Finally

definition *iEx* :: $iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Eventually

where *iEx* *I* *P* $\equiv \exists t \in I. P\ t$

syntax

-iAll :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ (\langle (3\Box - \cdot / -) \rangle [0, 0, 10] 10)$

-iEx :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ (\langle (3\Diamond - \cdot / -) \rangle [0, 0, 10] 10)$

syntax-consts

-iAll $\equiv iAll$ **and**

-iEx $\equiv iEx$

translations

$\Box\ t\ I. P \equiv CONST\ iAll\ I\ (\lambda t. P)$

$\Diamond\ t\ I. P \equiv CONST\ iEx\ I\ (\lambda t. P)$

Future temporal operator: Next

definition *iNext* :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Next

where *iNext* *t0* *I* *P* $\equiv P\ (inext\ t0\ I)$

Past temporal operator: Last/Previous

definition $iLast :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Last
where $iLast\ t0\ I\ P \equiv P\ (iprev\ t0\ I)$

syntax

$-iNext :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ (\langle (3 \circ - - - / -) \rangle [0, 0, 10] 10)$

$-iLast :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ (\langle (3 \ominus - - - / -) \rangle [0, 0, 10] 10)$

syntax-consts

$-iNext \equiv iNext$ **and**

$-iLast \equiv iLast$

translations

$\circ\ t\ t0\ I.\ P \equiv CONST\ iNext\ t0\ I\ (\lambda t.\ P)$

$\ominus\ t\ t0\ I.\ P \equiv CONST\ iLast\ t0\ I\ (\lambda t.\ P)$

lemma $\circ\ t\ 10\ [0\dots].\ (t + 10 > 10)$

<proof>

The following versions of Next and Last operator differ in the cases where no next/previous element exists or specified time point is not in interval: the weak versions return *True* and the strong versions return *False*.

definition $iNextWeak :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Weak Next
where $iNextWeak\ t0\ I\ P \equiv (\Box\ t\ \{inext\ t0\ I\} \downarrow >\ t0.\ P\ t)$

definition $iNextStrong :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Strong Next
where $iNextStrong\ t0\ I\ P \equiv (\Diamond\ t\ \{inext\ t0\ I\} \downarrow >\ t0.\ P\ t)$

definition $iLastWeak :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Weak Last
where $iLastWeak\ t0\ I\ P \equiv (\Box\ t\ \{iprev\ t0\ I\} \downarrow <\ t0.\ P\ t)$

definition $iLastStrong :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Strong Last
where $iLastStrong\ t0\ I\ P \equiv (\Diamond\ t\ \{iprev\ t0\ I\} \downarrow <\ t0.\ P\ t)$

syntax

$-iNextWeak :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ (\langle (3 \circ_W - - - / -) \rangle [0, 0, 10] 10)$

$-iNextStrong :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ (\langle (3 \circ_S - - - / -) \rangle [0, 0, 10] 10)$

$-iLastWeak :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ (\langle (3 \ominus_W - - - / -) \rangle [0, 0, 10] 10)$

$-iLastStrong :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ (\langle (3 \ominus_S - - - / -) \rangle [0, 0, 10] 10)$

syntax-consts

$-iNextWeak \equiv iNextWeak$ **and**

$-iNextStrong \equiv iNextStrong$ **and**

$-iLastWeak \equiv iLastWeak$ **and**

$-iLastStrong \equiv iLastStrong$

translations

$$\begin{aligned}
\circ_W t \ t0 \ I. \ P &\equiv \text{CONST } iNextWeak \ t0 \ I \ (\lambda t. \ P) \\
\circ_S t \ t0 \ I. \ P &\equiv \text{CONST } iNextStrong \ t0 \ I \ (\lambda t. \ P) \\
\ominus_W t \ t0 \ I. \ P &\equiv \text{CONST } iLastWeak \ t0 \ I \ (\lambda t. \ P) \\
\ominus_S t \ t0 \ I. \ P &\equiv \text{CONST } iLastStrong \ t0 \ I \ (\lambda t. \ P)
\end{aligned}$$

Some examples for Next and Last operator

lemma $\circ \ t \ 5 \ [0 \dots, 10]. \ ([0::int, 10, 20, 30, 40, 50, 60, 70, 80, 90] \ ! \ t < 80)$
<proof>

lemma $\ominus \ t \ 5 \ [0 \dots, 10]. \ ([0::int, 10, 20, 30, 40, 50, 60, 70, 80, 90] \ ! \ t < 80)$
<proof>

Temporal Until operator

definition $iUntil :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Until
where $iUntil \ I \ P \ Q \equiv \Diamond \ t \ I. \ Q \ t \wedge (\Box \ t' (I \downarrow < t). \ P \ t')$

Temporal Since operator (past operator corresponding to Until)

definition $iSince :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Since
where $iSince \ I \ P \ Q \equiv \Diamond \ t \ I. \ Q \ t \wedge (\Box \ t' (I \downarrow > t). \ P \ t')$

syntax

$$\begin{aligned}
-iUntil &:: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \\
&(\langle \langle \cdot / - \ (3U \ - \ -) / - \rangle \rangle [10, 0, 0, 0, 10] \ 10) \\
-iSince &:: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \\
&(\langle \langle \cdot / - \ (3S \ - \ -) / - \rangle \rangle [10, 0, 0, 0, 10] \ 10)
\end{aligned}$$

syntax-consts

$$\begin{aligned}
-iUntil &\equiv iUntil \text{ and} \\
-iSince &\equiv iSince
\end{aligned}$$

translations

$$\begin{aligned}
P. \ t \ U \ t' \ I. \ Q &\equiv \text{CONST } iUntil \ I \ (\lambda t. \ P) \ (\lambda t'. \ Q) \\
P. \ t \ S \ t' \ I. \ Q &\equiv \text{CONST } iSince \ I \ (\lambda t. \ P) \ (\lambda t'. \ Q)
\end{aligned}$$

definition $iWeakUntil :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ —
Weak Until/Wating-for/Unless

$$\textbf{where } iWeakUntil \ I \ P \ Q \equiv (\Box \ t \ I. \ P \ t) \vee (\Diamond \ t \ I. \ Q \ t \wedge (\Box \ t' (I \downarrow < t). \ P \ t'))$$

definition $iWeakSince :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ —
Weak Since/Back-to

$$\textbf{where } iWeakSince \ I \ P \ Q \equiv (\Box \ t \ I. \ P \ t) \vee (\Diamond \ t \ I. \ Q \ t \wedge (\Box \ t' (I \downarrow > t). \ P \ t'))$$

syntax

$$\begin{aligned}
-iWeakUntil &:: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \\
&(\langle \langle \cdot / - \ (3W \ - \ -) / - \rangle \rangle [10, 0, 0, 0, 10] \ 10) \\
-iWeakSince &:: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \\
&(\langle \langle \cdot / - \ (3B \ - \ -) / - \rangle \rangle [10, 0, 0, 0, 10] \ 10)
\end{aligned}$$

syntax-consts

-*iWeakUntil* $\Rightarrow$ *iWeakUntil* **and**
 -*iWeakSince* $\Rightarrow$ *iWeakSince*

translations

$P. t \mathcal{W} t' I. Q \Rightarrow \text{CONST } i\text{WeakUntil } I (\lambda t. P) (\lambda t'. Q)$
 $P. t \mathcal{B} t' I. Q \Rightarrow \text{CONST } i\text{WeakSince } I (\lambda t. P) (\lambda t'. Q)$

definition *iRelease* $:: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ —
 Release

where *iRelease* $I P Q \equiv (\Box t I. Q t) \vee (\Diamond t I. P t \wedge (\Box t' (I \downarrow \leq t). Q t'))$

definition *iTrigger* $:: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ —
 Trigger

where *iTrigger* $I P Q \equiv (\Box t I. Q t) \vee (\Diamond t I. P t \wedge (\Box t' (I \downarrow \geq t). Q t'))$

syntax

-*iRelease* $:: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $(\langle \langle \cdot / - (3\mathcal{R} - \cdot) / - \rangle \rangle [10, 0, 0, 0, 10] 10)$
 -*iTrigger* $:: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $(\langle \langle \cdot / - (3\mathcal{T} - \cdot) / - \rangle \rangle [10, 0, 0, 0, 10] 10)$

syntax-consts

-*iRelease* $\Rightarrow$ *iRelease* **and**
 -*iTrigger* $\Rightarrow$ *iTrigger*

translations

$P. t \mathcal{R} t' I. Q \Rightarrow \text{CONST } i\text{Release } I (\lambda t. P) (\lambda t'. Q)$
 $P. t \mathcal{T} t' I. Q \Rightarrow \text{CONST } i\text{Trigger } I (\lambda t. P) (\lambda t'. Q)$

lemmas *iTL-Next-defs* =

iNext-def *iLast-def*
iNextWeak-def *iLastWeak-def*
iNextStrong-def *iLastStrong-def*

lemmas *iTL-defs* =

iAll-def *iEx-def*
iUntil-def *iSince-def*
iWeakUntil-def *iWeakSince-def*
iRelease-def *iTrigger-def*
iTL-Next-defs

$\langle ML \rangle$

3.2 Basic lemmata for temporal operators

3.2.1 Intro/elim rules

lemma

iexI[intro]: $\llbracket P t; t \in I \rrbracket \Longrightarrow \Diamond t I. P t$ **and**
rev-iexI[intro?]: $\llbracket t \in I; P t \rrbracket \Longrightarrow \Diamond t I. P t$ **and**
iexE[elim!]: $\llbracket \Diamond t I. P t; \bigwedge t. \llbracket t \in I; P t \rrbracket \Longrightarrow Q \rrbracket \Longrightarrow Q$

$\langle proof \rangle$

lemma

iallI[intro]: $(\bigwedge t. t \in I \implies P\ t) \implies \Box\ t\ I. P\ t$ **and**
ispec[dest]: $\llbracket \Box\ t\ I. P\ t; t \in I \rrbracket \implies P\ t$ **and**
iallE[elim]: $\llbracket \Box\ t\ I. P\ t; P\ t \implies Q; t \notin I \implies Q \rrbracket \implies Q$
 $\langle proof \rangle$

lemma

iuntilI[intro]:
 $\llbracket Q\ t; (\bigwedge t'. t' \in I \downarrow < t \implies P\ t'); t \in I \rrbracket \implies P\ t'. t' \mathcal{U}\ t\ I. Q\ t$ **and**
rev-iuntilI[intro]:
 $\llbracket t \in I; Q\ t; (\bigwedge t'. t' \in I \downarrow < t \implies P\ t') \rrbracket \implies P\ t'. t' \mathcal{U}\ t\ I. Q\ t$
 $\langle proof \rangle$

lemma

iuntilE[elim]:
 $\llbracket P'\ t'. t' \mathcal{U}\ t\ I. P\ t; \bigwedge t. \llbracket t \in I; P\ t \rrbracket \implies Q \rrbracket \implies Q$
 $\langle proof \rangle$

lemma

isinceI[intro]:
 $\llbracket Q\ t; (\bigwedge t'. t' \in I \downarrow > t \implies P\ t'); t \in I \rrbracket \implies P\ t'. t' \mathcal{S}\ t\ I. Q\ t$ **and**
rev-isinceI[intro]:
 $\llbracket t \in I; Q\ t; (\bigwedge t'. t' \in I \downarrow > t \implies P\ t') \rrbracket \implies P\ t'. t' \mathcal{S}\ t\ I. Q\ t$
 $\langle proof \rangle$

lemma

isinceE[elim]:
 $\llbracket P'\ t'. t' \mathcal{S}\ t\ I. P\ t; \bigwedge t. \llbracket t \in I; P\ t \rrbracket \implies Q \rrbracket \implies Q$
 $\langle proof \rangle$

3.2.2 Rewrite rules for trivial simplification

lemma *iall-triv*[simp]: $(\Box\ t\ I. P) = ((\exists t. t \in I) \longrightarrow P)$
 $\langle proof \rangle$

lemma *ies-triv*[simp]: $(\Diamond\ t\ I. P) = ((\exists t. t \in I) \wedge P)$
 $\langle proof \rangle$

lemma *ies-conjL1*:

$(\Diamond\ t1\ I1. (P1\ t1 \wedge (\Diamond\ t2\ I2. P2\ t1\ t2))) =$
 $(\Diamond\ t1\ I1. \Diamond\ t2\ I2. P1\ t1 \wedge P2\ t1\ t2)$
 $\langle proof \rangle$

lemma *ies-conjR1*:

$(\Diamond\ t1\ I1. ((\Diamond\ t2\ I2. P2\ t1\ t2) \wedge P1\ t1)) =$
 $(\Diamond\ t1\ I1. \Diamond\ t2\ I2. P2\ t1\ t2 \wedge P1\ t1)$
 $\langle proof \rangle$

lemma *iex-conjL2*:

$$\begin{aligned} & (\Diamond t1\ I1. (P1\ t1 \wedge (\Diamond t2\ (I2\ t1). P2\ t1\ t2))) = \\ & (\Diamond t1\ I1. \Diamond t2\ (I2\ t1). P1\ t1 \wedge P2\ t1\ t2) \\ & \langle proof \rangle \end{aligned}$$

lemma *iex-conjR2*:

$$\begin{aligned} & (\Diamond t1\ I1. ((\Diamond t2\ (I2\ t1). P2\ t1\ t2) \wedge P1\ t1)) = \\ & (\Diamond t1\ I1. \Diamond t2\ (I2\ t1). P2\ t1\ t2 \wedge P1\ t1) \\ & \langle proof \rangle \end{aligned}$$

lemma *iex-commute*:

$$\begin{aligned} & (\Diamond t1\ I1. \Diamond t2\ I2. P\ t1\ t2) = \\ & (\Diamond t2\ I2. \Diamond t1\ I1. P\ t1\ t2) \\ & \langle proof \rangle \end{aligned}$$

lemma *iall-conjL1*:

$$\begin{aligned} & I2 \neq \{\} \implies \\ & (\Box t1\ I1. (P1\ t1 \wedge (\Box t2\ I2. P2\ t1\ t2))) = \\ & (\Box t1\ I1. \Box t2\ I2. P1\ t1 \wedge P2\ t1\ t2) \\ & \langle proof \rangle \end{aligned}$$

lemma *iall-conjR1*:

$$\begin{aligned} & I2 \neq \{\} \implies \\ & (\Box t1\ I1. ((\Box t2\ I2. P2\ t1\ t2) \wedge P1\ t1)) = \\ & (\Box t1\ I1. \Box t2\ I2. P2\ t1\ t2 \wedge P1\ t1) \\ & \langle proof \rangle \end{aligned}$$

lemma *iall-conjL2*:

$$\begin{aligned} & \Box t1\ I1. I2\ t1 \neq \{\} \implies \\ & (\Box t1\ I1. (P1\ t1 \wedge (\Box t2\ (I2\ t1). P2\ t1\ t2))) = \\ & (\Box t1\ I1. \Box t2\ (I2\ t1). P1\ t1 \wedge P2\ t1\ t2) \\ & \langle proof \rangle \end{aligned}$$

lemma *iall-conjR2*:

$$\begin{aligned} & \Box t1\ I1. I2\ t1 \neq \{\} \implies \\ & (\Box t1\ I1. ((\Box t2\ (I2\ t1). P2\ t1\ t2) \wedge P1\ t1)) = \\ & (\Box t1\ I1. \Box t2\ (I2\ t1). P2\ t1\ t2 \wedge P1\ t1) \\ & \langle proof \rangle \end{aligned}$$

lemma *iall-commute*:

$$\begin{aligned} & (\Box t1\ I1. \Box t2\ I2. P\ t1\ t2) = \\ & (\Box t2\ I2. \Box t1\ I1. P\ t1\ t2) \\ & \langle proof \rangle \end{aligned}$$

lemma *iall-conj-distrib*:

$$\begin{aligned} & (\Box t\ I. P\ t \wedge Q\ t) = ((\Box t\ I. P\ t) \wedge (\Box t\ I. Q\ t)) \\ & \langle proof \rangle \end{aligned}$$

lemma *iex-disj-distrib*:

$$(\Diamond t I. P t \vee Q t) = ((\Diamond t I. P t) \vee (\Diamond t I. Q t))$$

<proof>

lemma *iall-conj-distrib2*:

$$(\Box t1 I1. \Box t2 (I2 t1). P t1 t2 \wedge Q t1 t2) =$$

$$((\Box t1 I1. \Box t2 (I2 t1). P t1 t2) \wedge (\Box t1 I1. \Box t2 (I2 t1). Q t1 t2))$$

<proof>

lemma *iox-disj-distrib2*:

$$(\Diamond t1 I1. \Diamond t2 (I2 t1). P t1 t2 \vee Q t1 t2) =$$

$$((\Diamond t1 I1. \Diamond t2 (I2 t1). P t1 t2) \vee (\Diamond t1 I1. \Diamond t2 (I2 t1). Q t1 t2))$$

<proof>

lemma *iUntil-disj-distrib*:

$$(P t1. t1 \mathcal{U} t2 I. (Q1 t2 \vee Q2 t2)) = ((P t1. t1 \mathcal{U} t2 I. Q1 t2) \vee (P t1. t1 \mathcal{U} t2 I. Q2 t2))$$

<proof>

lemma *iSince-disj-distrib*:

$$(P t1. t1 \mathcal{S} t2 I. (Q1 t2 \vee Q2 t2)) = ((P t1. t1 \mathcal{S} t2 I. Q1 t2) \vee (P t1. t1 \mathcal{S} t2 I. Q2 t2))$$

<proof>

lemma

$$iNext\text{-}iff: (\bigcirc t t0 I. P t) = (\Box t [\dots 0] \oplus (inext t0 I). P t) \text{ and}$$

$$iLast\text{-}iff: (\ominus t t0 I. P t) = (\Box t [\dots 0] \oplus (iprev t0 I). P t)$$

<proof>

lemma

$$iNext\text{-}iEx\text{-}iff: (\bigcirc t t0 I. P t) = (\Diamond t [\dots 0] \oplus (inext t0 I). P t) \text{ and}$$

$$iLast\text{-}iEx\text{-}iff: (\ominus t t0 I. P t) = (\Diamond t [\dots 0] \oplus (iprev t0 I). P t)$$

<proof>

lemma *inext-singleton-cut-greater-not-empty-iff*:

$$(\{inext t0 I\} \Downarrow > t0 \neq \{\}) = (I \Downarrow > t0 \neq \{\} \wedge t0 \in I)$$

<proof>

lemma *iprev-singleton-cut-less-not-empty-iff*:

$$(\{iprev t0 I\} \Downarrow < t0 \neq \{\}) = (I \Downarrow < t0 \neq \{\} \wedge t0 \in I)$$

<proof>

lemma *inext-singleton-cut-greater-empty-iff*:

$$(\{inext t0 I\} \Downarrow > t0 = \{\}) = (I \Downarrow > t0 = \{\} \vee t0 \notin I)$$

<proof>

lemma *iprev-singleton-cut-less-empty-iff*:

$$(\{iprev t0 I\} \Downarrow < t0 = \{\}) = (I \Downarrow < t0 = \{\} \vee t0 \notin I)$$

<proof>

lemma *iNextWeak-iff* : $(\circ_W t \ t0 \ I. \ P \ t) = ((\circ \ t \ t0 \ I. \ P \ t) \vee (I \downarrow > t0 = \{\})) \vee t0 \notin I)$
 $\langle proof \rangle$

lemma *iNextStrong-iff* : $(\circ_S t \ t0 \ I. \ P \ t) = ((\circ \ t \ t0 \ I. \ P \ t) \wedge (I \downarrow > t0 \neq \{\})) \wedge t0 \in I)$
 $\langle proof \rangle$

lemma *iLastWeak-iff* : $(\ominus_W t \ t0 \ I. \ P \ t) = ((\ominus \ t \ t0 \ I. \ P \ t) \vee (I \downarrow < t0 = \{\})) \vee t0 \notin I)$
 $\langle proof \rangle$

lemma *iLastStrong-iff* : $(\ominus_S t \ t0 \ I. \ P \ t) = ((\ominus \ t \ t0 \ I. \ P \ t) \wedge (I \downarrow < t0 \neq \{\})) \wedge t0 \in I)$
 $\langle proof \rangle$

lemmas *iTL-Next-iff* =
iNext-iff *iLast-iff*
iNextWeak-iff *iNextStrong-iff*
iLastWeak-iff *iLastStrong-iff*

lemma
iNext-iff-singleton : $(\circ \ t \ t0 \ I. \ P \ t) = (\Box \ t \ \{\text{iNext } t0 \ I\}. \ P \ t)$ **and**
iLast-iff-singleton : $(\ominus \ t \ t0 \ I. \ P \ t) = (\Box \ t \ \{\text{iPrev } t0 \ I\}. \ P \ t)$
 $\langle proof \rangle$

lemmas *iNextLast-iff-singleton* = *iNext-iff-singleton* *iLast-iff-singleton*

lemma
iNext-iEx-iff-singleton : $(\circ \ t \ t0 \ I. \ P \ t) = (\Diamond \ t \ \{\text{iNext } t0 \ I\}. \ P \ t)$ **and**
iLast-iEx-iff-singleton : $(\ominus \ t \ t0 \ I. \ P \ t) = (\Diamond \ t \ \{\text{iPrev } t0 \ I\}. \ P \ t)$
 $\langle proof \rangle$

lemma
iNextWeak-iAll-conv: $(\circ_W t \ t0 \ I. \ P \ t) = (\Box \ t \ (\{\text{iNext } t0 \ I\} \downarrow > t0). \ P \ t)$ **and**
iNextStrong-iEx-conv: $(\circ_S t \ t0 \ I. \ P \ t) = (\Diamond \ t \ (\{\text{iNext } t0 \ I\} \downarrow > t0). \ P \ t)$ **and**
iLastWeak-iAll-conv: $(\ominus_W t \ t0 \ I. \ P \ t) = (\Box \ t \ (\{\text{iPrev } t0 \ I\} \downarrow < t0). \ P \ t)$ **and**
iLastStrong-iEx-conv: $(\ominus_S t \ t0 \ I. \ P \ t) = (\Diamond \ t \ (\{\text{iPrev } t0 \ I\} \downarrow < t0). \ P \ t)$
 $\langle proof \rangle$

lemma
iAll-True[simp]: $\Box \ t \ I. \ \text{True}$ **and**
iAll-False[simp]: $(\Box \ t \ I. \ \text{False}) = (I = \{\})$ **and**
iEx-True[simp]: $(\Diamond \ t \ I. \ \text{True}) = (I \neq \{\})$ **and**
iEx-False[simp]: $\neg (\Diamond \ t \ I. \ \text{False})$
 $\langle proof \rangle$

lemma *empty-iff-iAll-False*: $(I = \{\}) = (\Box t I. \text{False})$ *<proof>*

lemma *not-empty-iff-iEx-True*: $(I \neq \{\}) = (\Diamond t I. \text{True})$ *<proof>*

lemma

iNext-True: $\bigcirc t t0 I. \text{True}$ **and**

iNextWeak-True: $(\bigcirc_W t t0 I. \text{True})$ **and**

iNext-False: $\neg (\bigcirc t t0 I. \text{False})$ **and**

iNextStrong-False: $\neg (\bigcirc_S t t0 I. \text{False})$

<proof>

lemma

iNextStrong-True: $(\bigcirc_S t t0 I. \text{True}) = (I \Downarrow> t0 \neq \{\} \wedge t0 \in I)$ **and**

iNextWeak-False: $(\neg (\bigcirc_W t t0 I. \text{False})) = (I \Downarrow> t0 \neq \{\} \wedge t0 \in I)$

<proof>

lemma

iLast-True: $\ominus t t0 I. \text{True}$ **and**

iLastWeak-True: $(\ominus_W t t0 I. \text{True})$ **and**

iLast-False: $\neg (\ominus t t0 I. \text{False})$ **and**

iLastStrong-False: $\neg (\ominus_S t t0 I. \text{False})$

<proof>

lemma

iLastStrong-True: $(\ominus_S t t0 I. \text{True}) = (I \Downarrow< t0 \neq \{\} \wedge t0 \in I)$ **and**

iLastWeak-False: $(\neg (\ominus_W t t0 I. \text{False})) = (I \Downarrow< t0 \neq \{\} \wedge t0 \in I)$

<proof>

lemma *iUntil-True-left[simp]*: $(\text{True}. t' \mathcal{U} t I. Q t) = (\Diamond t I. Q t)$

<proof>

lemma *iUntil-True[simp]*: $(P t'. t' \mathcal{U} t I. \text{True}) = (I \neq \{\})$

<proof>

lemma *iUntil-False-left[simp]*: $(\text{False}. t' \mathcal{U} t I. Q t) = (I \neq \{\} \wedge Q (iMin I))$

<proof>

lemma *iUntil-False[simp]*: $\neg (P t'. t' \mathcal{U} t I. \text{False})$

<proof>

lemma *iSince-True-left[simp]*: $(\text{True}. t' \mathcal{S} t I. Q t) = (\Diamond t I. Q t)$

<proof>

lemma *iSince-True-if*:

$(P t'. t' \mathcal{S} t I. \text{True}) =$

$(\text{if finite } I \text{ then } I \neq \{\} \text{ else } \Diamond t1 I. \Box t2 (I \Downarrow> t1). P t2)$

<proof>

corollary *iSince-True-finite[simp]*: $\text{finite } I \implies (P \ t'. \ t' \mathcal{S} \ t \ I. \ \text{True}) = (I \neq \{\})$
 $\langle \text{proof} \rangle$

lemma *iSince-False-left[simp]*: $(\text{False}. \ t' \mathcal{S} \ t \ I. \ Q \ t) = (\text{finite } I \wedge I \neq \{\} \wedge Q \ (Max \ I))$
 $\langle \text{proof} \rangle$

lemma *iSince-False[simp]*: $\neg (P \ t'. \ t' \mathcal{S} \ t \ I. \ \text{False})$
 $\langle \text{proof} \rangle$

lemma *iWeakUntil-True-left[simp]*: $\text{True}. \ t' \mathcal{W} \ t \ I. \ Q \ t$
 $\langle \text{proof} \rangle$

lemma *iWeakUntil-True[simp]*: $P \ t'. \ t' \mathcal{W} \ t \ I. \ \text{True}$
 $\langle \text{proof} \rangle$

lemma *iWeakUntil-False-left[simp]*: $(\text{False}. \ t' \mathcal{W} \ t \ I. \ Q \ t) = (I = \{\} \vee Q \ (iMin \ I))$
 $\langle \text{proof} \rangle$

lemma *iWeakUntil-False[simp]*: $(P \ t'. \ t' \mathcal{W} \ t \ I. \ \text{False}) = (\Box \ t \ I. \ P \ t)$
 $\langle \text{proof} \rangle$

lemma *iWeakSince-True-left[simp]*: $\text{True}. \ t' \mathcal{B} \ t \ I. \ Q \ t$
 $\langle \text{proof} \rangle$

lemma *iWeakSince-True-disj*:
 $(P \ t'. \ t' \mathcal{B} \ t \ I. \ \text{True}) =$
 $(I = \{\} \vee (\Diamond \ t1 \ I. \ \Box \ t2 \ (I \Downarrow t1). \ P \ t2))$
 $\langle \text{proof} \rangle$

lemma *iWeakSince-True-finite[simp]*: $\text{finite } I \implies (P \ t'. \ t' \mathcal{B} \ t \ I. \ \text{True})$
 $\langle \text{proof} \rangle$

lemma *iWeakSince-False-left[simp]*: $(\text{False}. \ t' \mathcal{B} \ t \ I. \ Q \ t) = (I = \{\} \vee (\text{finite } I \wedge Q \ (Max \ I)))$
 $\langle \text{proof} \rangle$

lemma *iWeakSince-False[simp]*: $(P \ t'. \ t' \mathcal{B} \ t \ I. \ \text{False}) = (\Box \ t \ I. \ P \ t)$
 $\langle \text{proof} \rangle$

lemma *iRelease-True-left[simp]*: $(\text{True}. \ t' \mathcal{R} \ t \ I. \ Q \ t) = (I = \{\} \vee Q \ (iMin \ I))$
 $\langle \text{proof} \rangle$

lemma *iRelease-True[simp]*: $P \ t'. \ t' \mathcal{R} \ t \ I. \ \text{True}$
 $\langle \text{proof} \rangle$

lemma *iRelease-False-left[simp]*: $(\text{False}. \ t' \mathcal{R} \ t \ I. \ Q \ t) = (\Box \ t \ I. \ Q \ t)$
 $\langle \text{proof} \rangle$

lemma *iRelease-False[simp]*: $(P\ t'.\ t'\ \mathcal{R}\ t\ I.\ \text{False}) = (I = \{\})$
 $\langle \text{proof} \rangle$

lemma *iTrigger-True-left[simp]*: $(\text{True}.\ t'\ \mathcal{T}\ t\ I.\ Q\ t) = (I = \{\} \vee (\Diamond\ t1\ I.\ \Box\ t2\ (I \downarrow_{\geq} t1).\ Q\ t2))$
 $\langle \text{proof} \rangle$

lemma *iTrigger-True[simp]*: $P\ t'.\ t'\ \mathcal{T}\ t\ I.\ \text{True}$
 $\langle \text{proof} \rangle$

lemma *iTrigger-False-left[simp]*: $(\text{False}.\ t'\ \mathcal{T}\ t\ I.\ Q\ t) = (\Box\ t\ I.\ Q\ t)$
 $\langle \text{proof} \rangle$

lemma *iTrigger-False[simp]*: $(P\ t'.\ t'\ \mathcal{T}\ t\ I.\ \text{False}) = (I = \{\})$
 $\langle \text{proof} \rangle$

lemma
iUntil-TrueTrue[simp]: $(\text{True}.\ t'\ \mathcal{U}\ t\ I.\ \text{True}) = (I \neq \{\})$ **and**
iSince-TrueTrue[simp]: $(\text{True}.\ t'\ \mathcal{S}\ t\ I.\ \text{True}) = (I \neq \{\})$ **and**
iWeakUntil-TrueTrue[simp]: $\text{True}.\ t'\ \mathcal{W}\ t\ I.\ \text{True}$ **and**
iWeakSince-TrueTrue[simp]: $\text{True}.\ t'\ \mathcal{B}\ t\ I.\ \text{True}$ **and**
iRelease-TrueTrue[simp]: $\text{True}.\ t'\ \mathcal{R}\ t\ I.\ \text{True}$ **and**
iTrigger-TrueTrue[simp]: $\text{True}.\ t'\ \mathcal{T}\ t\ I.\ \text{True}$
 $\langle \text{proof} \rangle$

3.2.3 Empty sets and singletons

lemma *iAll-empty[simp]*: $\Box\ t\ \{\}.\ P\ t\ \langle \text{proof} \rangle$

lemma *iEx-empty[simp]*: $\neg (\Diamond\ t\ \{\}.\ P\ t)\ \langle \text{proof} \rangle$

lemma *iUntil-empty[simp]*: $\neg (P\ t0.\ t0\ \mathcal{U}\ t1\ \{\}.\ Q\ t1)\ \langle \text{proof} \rangle$

lemma *iSince-empty[simp]*: $\neg (P\ t0.\ t0\ \mathcal{S}\ t1\ \{\}.\ Q\ t1)\ \langle \text{proof} \rangle$

lemma *iWeakUntil-empty[simp]*: $P\ t0.\ t0\ \mathcal{W}\ t1\ \{\}.\ Q\ t1\ \langle \text{proof} \rangle$

lemma *iWeakSince-empty[simp]*: $P\ t0.\ t0\ \mathcal{B}\ t1\ \{\}.\ Q\ t1\ \langle \text{proof} \rangle$

lemma *iRelease-empty[simp]*: $P\ t0.\ t0\ \mathcal{R}\ t1\ \{\}.\ Q\ t1\ \langle \text{proof} \rangle$

lemma *iTrigger-empty[simp]*: $P\ t0.\ t0\ \mathcal{T}\ t1\ \{\}.\ Q\ t1\ \langle \text{proof} \rangle$

lemmas *iTL-empty* =
iAll-empty *iEx-empty*
iUntil-empty *iSince-empty*
iWeakUntil-empty *iWeakSince-empty*
iRelease-empty *iTrigger-empty*

lemma *iAll-singleton[simp]*: $(\Box\ t'\ \{t\}.\ P\ t') = P\ t\ \langle \text{proof} \rangle$

lemma *iEx-singleton[simp]*: $(\Diamond\ t'\ \{t\}.\ P\ t') = P\ t\ \langle \text{proof} \rangle$

lemma *iUntil-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{U}\ t1\ \{t\}.\ Q\ t1) = Q\ t$

$\langle \text{proof} \rangle$

lemma *iSince-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{S}\ t1\ \{t\}.\ Q\ t1) = Q\ t$
 $\langle \text{proof} \rangle$

lemma *iWeakUntil-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{W}\ t1\ \{t\}.\ Q\ t1) = (P\ t \vee Q\ t)$
 $\langle \text{proof} \rangle$

lemma *iWeakSince-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{B}\ t1\ \{t\}.\ Q\ t1) = (P\ t \vee Q\ t)$
 $\langle \text{proof} \rangle$

lemma *iRelease-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{R}\ t1\ \{t\}.\ Q\ t1) = Q\ t$
 $\langle \text{proof} \rangle$

lemma *iTrigger-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{T}\ t1\ \{t\}.\ Q\ t1) = Q\ t$
 $\langle \text{proof} \rangle$

lemmas *iTL-singleton* =
iAll-singleton iEx-singleton
iUntil-singleton iSince-singleton
iWeakUntil-singleton iWeakSince-singleton
iRelease-singleton iTrigger-singleton

3.2.4 Conversions between temporal operators

lemma *iAll-iEx-conv*: $(\Box\ t\ I.\ P\ t) = (\neg (\Diamond\ t\ I.\ \neg P\ t))$ $\langle \text{proof} \rangle$

lemma *iEx-iAll-conv*: $(\Diamond\ t\ I.\ P\ t) = (\neg (\Box\ t\ I.\ \neg P\ t))$ $\langle \text{proof} \rangle$

lemma *not-iAll[simp]*: $(\neg (\Box\ t\ I.\ P\ t)) = (\Diamond\ t\ I.\ \neg P\ t)$ $\langle \text{proof} \rangle$

lemma *not-iEx[simp]*: $(\neg (\Diamond\ t\ I.\ P\ t)) = (\Box\ t\ I.\ \neg P\ t)$ $\langle \text{proof} \rangle$

lemma *iUntil-iEx-conv*: $(\text{True}.\ t'\ \mathcal{U}\ t\ I.\ P\ t) = (\Diamond\ t\ I.\ P\ t)$ $\langle \text{proof} \rangle$

lemma *iSince-iEx-conv*: $(\text{True}.\ t'\ \mathcal{S}\ t\ I.\ P\ t) = (\Diamond\ t\ I.\ P\ t)$ $\langle \text{proof} \rangle$

lemma *iRelease-iAll-conv*: $(\text{False}.\ t'\ \mathcal{R}\ t\ I.\ P\ t) = (\Box\ t\ I.\ P\ t)$
 $\langle \text{proof} \rangle$

lemma *iTrigger-iAll-conv*: $(\text{False}.\ t'\ \mathcal{T}\ t\ I.\ P\ t) = (\Box\ t\ I.\ P\ t)$
 $\langle \text{proof} \rangle$

lemma *iWeakUntil-iUntil-conv*: $(P\ t'.\ t'\ \mathcal{W}\ t\ I.\ Q\ t) = ((P\ t'.\ t'\ \mathcal{U}\ t\ I.\ Q\ t) \vee (\Box\ t\ I.\ P\ t))$
 $\langle \text{proof} \rangle$

lemma *iWeakSince-iSince-conv*: $(P\ t'.\ t'\ \mathcal{B}\ t\ I.\ Q\ t) = ((P\ t'.\ t'\ \mathcal{S}\ t\ I.\ Q\ t) \vee (\Box\ t\ I.\ P\ t))$
 $\langle \text{proof} \rangle$

lemma *iUntil-iWeakUntil-conv*: $(P\ t'.\ t'\ \mathcal{U}\ t\ I.\ Q\ t) = ((P\ t'.\ t'\ \mathcal{W}\ t\ I.\ Q\ t) \wedge (\Diamond\ t\ I.\ P\ t))$

$t I. Q t))$
 $\langle proof \rangle$

lemma *iSince-iWeakSince-conv*: $(P t'. t' \mathcal{S} t I. Q t) = ((P t'. t' \mathcal{B} t I. Q t) \wedge (\Diamond t I. Q t))$
 $\langle proof \rangle$

lemma *iRelease-iWeakUntil-conv*: $(P t'. t' \mathcal{R} t I. Q t) = (Q t'. t' \mathcal{W} t I. (Q t \wedge P t))$
 $\langle proof \rangle$

lemma *iRelease-iUntil-conv*: $(P t'. t' \mathcal{R} t I. Q t) = ((\Box t I. Q t) \vee (Q t'. t' \mathcal{U} t I. (Q t \wedge P t)))$
 $\langle proof \rangle$

lemma *iTrigger-iWeakSince-conv*: $(P t'. t' \mathcal{T} t I. Q t) = (Q t'. t' \mathcal{B} t I. (Q t \wedge P t))$
 $\langle proof \rangle$

lemma *iTrigger-iSince-conv*: $(P t'. t' \mathcal{T} t I. Q t) = ((\Box t I. Q t) \vee (Q t'. t' \mathcal{S} t I. (Q t \wedge P t)))$
 $\langle proof \rangle$

lemma *iRelease-not-iUntil-conv*: $(P t'. t' \mathcal{R} t I. Q t) = (\neg (\neg P t'. t' \mathcal{U} t I. \neg Q t))$
 $\langle proof \rangle$

lemma *iUntil-not-iRelease-conv*: $(P t'. t' \mathcal{U} t I. Q t) = (\neg (\neg P t'. t' \mathcal{R} t I. \neg Q t))$
 $\langle proof \rangle$

The Trigger operator $\mathcal{T}$ is a past operator, so that it is used for time intervals, that are bounded by a current time point, and thus are finite. For an infinite interval the stated relation to the Since operator $\mathcal{S}$ would not be fulfilled.

lemma *iTrigger-not-iSince-conv*: $finite I \implies (P t'. t' \mathcal{T} t I. Q t) = (\neg (\neg P t'. t' \mathcal{S} t I. \neg Q t))$
 $\langle proof \rangle$

lemma *iSince-not-iTrigger-conv*: $finite I \implies (P t'. t' \mathcal{S} t I. Q t) = (\neg (\neg P t'. t' \mathcal{T} t I. \neg Q t))$
 $\langle proof \rangle$

lemma *not-iUntil*:
 $(\neg (P t1. t1 \mathcal{U} t2 I. Q t2)) =$
 $(\Box t I. (Q t \longrightarrow (\Diamond t' (I \downarrow < t). \neg P t')))$
 $\langle proof \rangle$

lemma *not-iSince*:

$$\begin{aligned}
& (\neg (P \ t1. \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ t2)) = \\
& (\Box \ t \ I. (Q \ t \longrightarrow (\Diamond \ t' (I \downarrow > t). \neg P \ t'))) \\
\langle proof \rangle
\end{aligned}$$

lemma *iWeakUntil-conj-iUntil-conv:*

$$\begin{aligned}
& (P \ t1. \ t1 \ \mathcal{W} \ t2 \ I. (P \ t2 \wedge Q \ t2)) = (\neg (\neg Q \ t1. \ t1 \ \mathcal{U} \ t2 \ I. \neg P \ t2)) \\
\langle proof \rangle
\end{aligned}$$

lemma *iUntil-disj-iUntil-conv:*

$$\begin{aligned}
& (P \ t1 \vee Q \ t1. \ t1 \ \mathcal{U} \ t2 \ I. \ Q \ t2) = \\
& (P \ t1. \ t1 \ \mathcal{U} \ t2 \ I. \ Q \ t2) \\
\langle proof \rangle
\end{aligned}$$

lemma *iWeakUntil-disj-iWeakUntil-conv:*

$$\begin{aligned}
& (P \ t1 \vee Q \ t1. \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2) = \\
& (P \ t1. \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2) \\
\langle proof \rangle
\end{aligned}$$

lemma *iWeakUntil-iUntil-conj-conv:*

$$\begin{aligned}
& (P \ t1. \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2) = \\
& (\neg (\neg Q \ t1. \ t1 \ \mathcal{U} \ t2 \ I. (\neg P \ t2 \wedge \neg Q \ t2))) \\
\langle proof \rangle
\end{aligned}$$

Negation and temporal operators

lemma

$$\begin{aligned}
& not-iNext[simp]: (\neg (\bigcirc \ t \ t0 \ I. \ P \ t)) = (\bigcirc \ t \ t0 \ I. \neg P \ t) \text{ and} \\
& not-iNextWeak[simp]: (\neg (\bigcirc_W \ t \ t0 \ I. \ P \ t)) = (\bigcirc_S \ t \ t0 \ I. \neg P \ t) \text{ and} \\
& not-iNextStrong[simp]: (\neg (\bigcirc_S \ t \ t0 \ I. \ P \ t)) = (\bigcirc_W \ t \ t0 \ I. \neg P \ t) \text{ and} \\
& not-iLast[simp]: (\neg (\ominus \ t \ t0 \ I. \ P \ t)) = (\ominus \ t \ t0 \ I. \neg P \ t) \text{ and} \\
& not-iLastWeak[simp]: (\neg (\ominus_W \ t \ t0 \ I. \ P \ t)) = (\ominus_S \ t \ t0 \ I. \neg P \ t) \text{ and} \\
& not-iLastStrong[simp]: (\neg (\ominus_S \ t \ t0 \ I. \ P \ t)) = (\ominus_W \ t \ t0 \ I. \neg P \ t) \\
\langle proof \rangle
\end{aligned}$$

lemma *not-iWeakUntil:*

$$\begin{aligned}
& (\neg (P \ t1. \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2)) = \\
& (((\Box \ t \ I. (Q \ t \longrightarrow (\Diamond \ t' (I \downarrow < t). \neg P \ t'))) \wedge (\Diamond \ t \ I. \neg P \ t)) \\
\langle proof \rangle
\end{aligned}$$

lemma *not-iWeakSince:*

$$\begin{aligned}
& (\neg (P \ t1. \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ t2)) = \\
& (((\Box \ t \ I. (Q \ t \longrightarrow (\Diamond \ t' (I \downarrow > t). \neg P \ t'))) \wedge (\Diamond \ t \ I. \neg P \ t)) \\
\langle proof \rangle
\end{aligned}$$

lemma *not-iRelease:*

$$\begin{aligned}
& (\neg (P \ t'. \ t' \ \mathcal{R} \ t \ I. \ Q \ t)) = \\
& ((\Diamond \ t \ I. \neg Q \ t) \wedge (\Box \ t \ I. P \ t \longrightarrow (\Diamond \ t \ I \downarrow \leq t. \neg Q \ t))) \\
\langle proof \rangle
\end{aligned}$$

lemma *not-iRelease-iUntil-conv:*

$$(\neg (P \ t'. \ t' \ \mathcal{R} \ t \ I. \ Q \ t)) = (\neg P \ t'. \ t' \ \mathcal{U} \ t \ I. \neg Q \ t)$$

<proof>

lemma *not-iTrigger*:

$$(\neg (P \ t'. \ t' \ \mathcal{T} \ t \ I. \ Q \ t)) =$$

$$((\Diamond t \ I. \neg Q \ t) \wedge (\Box t \ I. \neg P \ t \vee (\Diamond t \ I \downarrow \geq t. \neg Q \ t)))$$

<proof>

lemma *not-iTrigger-iSince-conv*:

$$\text{finite } I \implies (\neg (P \ t'. \ t' \ \mathcal{T} \ t \ I. \ Q \ t)) = (\neg P \ t'. \ t' \ \mathcal{S} \ t \ I. \neg Q \ t)$$

<proof>

3.2.5 Some implication results

lemma *all-imp-iall*: $\forall x. P \ x \implies \Box t \ I. P \ t$ *<proof>*

lemma *bex-imp-lex*: $\Diamond t \ I. P \ t \implies \exists x. P \ x$ *<proof>*

lemma *iAll-imp-iEx*: $I \neq \{\} \implies \Box t \ I. P \ t \implies \Diamond t \ I. P \ t$ *<proof>*

lemma *i-set-iAll-imp-iEx*: $I \in \text{i-set} \implies \Box t \ I. P \ t \implies \Diamond t \ I. P \ t$

<proof>

lemmas *iT-iAll-imp-iEx* = *iT-not-empty*[*THEN* *iAll-imp-iEx*]

lemma *iUntil-imp-iEx*: $P \ t1. \ t1 \ \mathcal{U} \ t2 \ I. \ Q \ t2 \implies \Diamond t \ I. \ Q \ t$

<proof>

lemma *iSince-imp-iEx*: $P \ t1. \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ t2 \implies \Diamond t \ I. \ Q \ t$

<proof>

lemma *iall-subset-imp-iall*: $\llbracket \Box t \ B. \ P \ t; A \subseteq B \rrbracket \implies \Box t \ A. \ P \ t$

<proof>

lemma *iex-subset-imp-iex*: $\llbracket \Diamond t \ A. \ P \ t; A \subseteq B \rrbracket \implies \Diamond t \ B. \ P \ t$

<proof>

lemma *iall-mp*: $\llbracket \Box t \ I. \ P \ t \longrightarrow Q \ t; \Box t \ I. \ P \ t \rrbracket \implies \Box t \ I. \ Q \ t$ *<proof>*

lemma *iex-mp*: $\llbracket \Box t \ I. \ P \ t \longrightarrow Q \ t; \Diamond t \ I. \ P \ t \rrbracket \implies \Diamond t \ I. \ Q \ t$ *<proof>*

lemma *iUntil-imp*:

$$\llbracket P1 \ t1. \ t1 \ \mathcal{U} \ t2 \ I. \ Q \ t2; \Box t \ I. \ P1 \ t \longrightarrow P2 \ t \rrbracket \implies P2 \ t1. \ t1 \ \mathcal{U} \ t2 \ I. \ Q \ t2$$

<proof>

lemma *iSince-imp*:

$$\llbracket P1 \ t1. \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ t2; \Box t \ I. \ P1 \ t \longrightarrow P2 \ t \rrbracket \implies P2 \ t1. \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ t2$$

<proof>

lemma *iWeakUntil-imp*:

$$\llbracket P1 \ t1. \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2; \Box t \ I. \ P1 \ t \longrightarrow P2 \ t \rrbracket \implies P2 \ t1. \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2$$

$\langle proof \rangle$

lemma *iWeakSince-imp*:

$\llbracket P1\ t1.\ t1\ \mathcal{B}\ t2\ I.\ Q\ t2; \Box\ t\ I.\ P1\ t \longrightarrow P2\ t \rrbracket \Longrightarrow P2\ t1.\ t1\ \mathcal{B}\ t2\ I.\ Q\ t2$
 $\langle proof \rangle$

lemma *iRelease-imp*:

$\llbracket P1\ t1.\ t1\ \mathcal{R}\ t2\ I.\ Q\ t2; \Box\ t\ I.\ P1\ t \longrightarrow P2\ t \rrbracket \Longrightarrow P2\ t1.\ t1\ \mathcal{R}\ t2\ I.\ Q\ t2$
 $\langle proof \rangle$

lemma *iTrigger-imp*:

$\llbracket P1\ t1.\ t1\ \mathcal{T}\ t2\ I.\ Q\ t2; \Box\ t\ I.\ P1\ t \longrightarrow P2\ t \rrbracket \Longrightarrow P2\ t1.\ t1\ \mathcal{T}\ t2\ I.\ Q\ t2$
 $\langle proof \rangle$

lemma *iMin-imp-iUntil*:

$\llbracket I \neq \{\}; Q\ (iMin\ I) \rrbracket \Longrightarrow P\ t'.\ t'\ \mathcal{U}\ t\ I.\ Q\ t$
 $\langle proof \rangle$

lemma *Max-imp-iSince*:

$\llbracket finite\ I; I \neq \{\}; Q\ (Max\ I) \rrbracket \Longrightarrow P\ t'.\ t'\ \mathcal{S}\ t\ I.\ Q\ t$
 $\langle proof \rangle$

3.2.6 Congruence rules for temporal operators' predicates

lemma *iAll-cong*: $\Box\ t\ I.\ f\ t = g\ t \Longrightarrow (\Box\ t\ I.\ P\ (f\ t)\ t) = (\Box\ t\ I.\ P\ (g\ t)\ t)$
 $\langle proof \rangle$

lemma *iEx-cong*: $\Box\ t\ I.\ f\ t = g\ t \Longrightarrow (\Diamond\ t\ I.\ P\ (f\ t)\ t) = (\Diamond\ t\ I.\ P\ (g\ t)\ t)$
 $\langle proof \rangle$

lemma *iUntil-cong1*:

$\Box\ t\ I.\ f\ t = g\ t \Longrightarrow$
 $(P\ (f\ t1)\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ t2) = (P\ (g\ t1)\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ t2)$
 $\langle proof \rangle$

lemma *iUntil-cong2*:

$\Box\ t\ I.\ f\ t = g\ t \Longrightarrow$
 $(P\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ (f\ t2)\ t2) = (P\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ (g\ t2)\ t2)$
 $\langle proof \rangle$

lemma *iSince-cong1*:

$\Box\ t\ I.\ f\ t = g\ t \Longrightarrow$
 $(P\ (f\ t1)\ t1.\ t1\ \mathcal{S}\ t2\ I.\ Q\ t2) = (P\ (g\ t1)\ t1.\ t1\ \mathcal{S}\ t2\ I.\ Q\ t2)$
 $\langle proof \rangle$

lemma *iSince-cong2*:

$\Box\ t\ I.\ f\ t = g\ t \Longrightarrow$
 $(P\ t1.\ t1\ \mathcal{S}\ t2\ I.\ Q\ (f\ t2)\ t2) = (P\ t1.\ t1\ \mathcal{S}\ t2\ I.\ Q\ (g\ t2)\ t2)$

$\langle proof \rangle$

lemma *bex-subst*:

$$\begin{aligned} & \forall x \in A. P x \longrightarrow (Q x = Q' x) \implies \\ & (\exists x \in A. P x \wedge Q x) = (\exists x \in A. P x \wedge Q' x) \end{aligned}$$

$\langle proof \rangle$

lemma *iEx-subst*:

$$\begin{aligned} & \Box t I. P t \longrightarrow (Q t = Q' t) \implies \\ & (\Diamond t I. P t \wedge Q t) = (\Diamond t I. P t \wedge Q' t) \end{aligned}$$

$\langle proof \rangle$

3.2.7 Temporal operators with set unions/intersections and subsets

lemma *iAll-subset*: $\llbracket A \subseteq B; \Box t B. P t \rrbracket \implies \Box t A. P t$
 $\langle proof \rangle$

lemma *iEx-subset*: $\llbracket A \subseteq B; \Diamond t A. P t \rrbracket \implies \Diamond t B. P t$
 $\langle proof \rangle$

lemma *iUntil-append*:

$$\begin{aligned} & \llbracket \text{finite } A; \text{Max } A \leq i\text{Min } B \rrbracket \implies \\ & P t1. t1 \mathcal{U} t2 A. Q t2 \implies P t1. t1 \mathcal{U} t2 (A \cup B). Q t2 \end{aligned}$$

$\langle proof \rangle$

lemma *iSince-prepend*:

$$\begin{aligned} & \llbracket \text{finite } A; \text{Max } A \leq i\text{Min } B \rrbracket \implies \\ & P t1. t1 \mathcal{S} t2 B. Q t2 \implies P t1. t1 \mathcal{S} t2 (A \cup B). Q t2 \end{aligned}$$

$\langle proof \rangle$

lemma

$$\begin{aligned} & \textit{iAll-union}: \llbracket \Box t A. P t; \Box t B. P t \rrbracket \implies \Box t (A \cup B). P t \textbf{ and } \\ & \textit{iAll-union-conv}: (\Box t A \cup B. P t) = ((\Box t A. P t) \wedge (\Box t B. P t)) \end{aligned}$$

$\langle proof \rangle$

lemma

$$\begin{aligned} & \textit{iEx-union}: (\Diamond t A. P t) \vee (\Diamond t B. P t) \implies \Diamond t (A \cup B). P t \textbf{ and } \\ & \textit{iEx-union-conv}: (\Diamond t (A \cup B). P t) = ((\Diamond t A. P t) \vee (\Diamond t B. P t)) \end{aligned}$$

$\langle proof \rangle$

lemma *iAll-inter*: $(\Box t A. P t) \vee (\Box t B. P t) \implies \Box t (A \cap B). P t$ $\langle proof \rangle$

lemma *not-iEx-inter*:

$$\exists A B P. (\Diamond t A. P t) \wedge (\Diamond t B. P t) \wedge \neg (\Diamond t (A \cap B). P t)$$

$\langle proof \rangle$

lemma

iAll-insert: $\llbracket P t; \Box t I. P t \rrbracket \implies \Box t' (\text{insert } t I). P t'$ **and**
iAll-insert-conv: $(\Box t' (\text{insert } t I). P t') = (P t \wedge (\Box t' I. P t'))$
 $\langle \text{proof} \rangle$

lemma

iEx-insert: $\llbracket P t \vee (\Diamond t I. P t) \rrbracket \implies \Diamond t' (\text{insert } t I). P t'$ **and**
iEx-insert-conv: $(\Diamond t' (\text{insert } t I). P t') = (P t \vee (\Diamond t' I. P t'))$
 $\langle \text{proof} \rangle$

3.3 Further results for temporal operators

lemma *Collect-minI-iEx*: $\Diamond t I. P t \implies \Diamond t I. P t \wedge (\Box t' (I \downarrow < t). \neg P t')$
 $\langle \text{proof} \rangle$

lemma *iUntil-disj-conv1*:

$I \neq \{\}$ $\implies$
 $(P t'. t' \mathcal{U} t I. Q t) = (Q (iMin I) \vee (P t'. t' \mathcal{U} t I. Q t \wedge iMin I < t))$
 $\langle \text{proof} \rangle$

lemma *iSince-disj-conv1*:

$\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies$
 $(P t'. t' \mathcal{S} t I. Q t) = (Q (Max I) \vee (P t'. t' \mathcal{S} t I. Q t \wedge t < Max I))$
 $\langle \text{proof} \rangle$

lemma *iUntil-next*:

$I \neq \{\}$ $\implies$
 $(P t'. t' \mathcal{U} t I. Q t) =$
 $(Q (iMin I) \vee (P (iMin I) \wedge (P t'. t' \mathcal{U} t (I \downarrow > (iMin I)). Q t)))$
 $\langle \text{proof} \rangle$

lemma *iSince-prev*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies$

$(P t'. t' \mathcal{S} t I. Q t) =$
 $(Q (Max I) \vee (P (Max I) \wedge (P t'. t' \mathcal{S} t (I \downarrow < Max I). Q t)))$
 $\langle \text{proof} \rangle$

lemma *iNext-induct-rule*:

$\llbracket P (iMin I); \Box t I. (P t \longrightarrow (\bigcirc t' t I. P t')); t \in I \rrbracket \implies P t$
 $\langle \text{proof} \rangle$

lemma *iNext-induct*:

$\llbracket P (iMin I); \Box t I. (P t \longrightarrow (\bigcirc t' t I. P t')) \rrbracket \implies \Box t I. P t$
 $\langle \text{proof} \rangle$

lemma *iLast-induct-rule*:

$\llbracket P (Max I); \Box t I. (P t \longrightarrow (\ominus t' t I. P t')); \text{finite } I; t \in I \rrbracket \implies P t$
 $\langle \text{proof} \rangle$

lemma *iLast-induct*:

$\llbracket P (Max I); \Box t I. (P t \longrightarrow (\ominus t' t I. P t')); \text{finite } I \rrbracket \implies \Box t I. P t$

$\langle \text{proof} \rangle$

lemma *iUntil-conj-not*: $((P \ t1 \wedge \neg Q \ t1). \ t1 \ \mathcal{U} \ t2 \ I. \ Q \ t2) = (P \ t1. \ t1 \ \mathcal{U} \ t2 \ I. \ Q \ t2)$

$\langle \text{proof} \rangle$

lemma *iWeakUntil-conj-not*: $((P \ t1 \wedge \neg Q \ t1). \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2) = (P \ t1. \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2)$

$\langle \text{proof} \rangle$

lemma *iSince-conj-not*: $\text{finite } I \implies$

$((P \ t1 \wedge \neg Q \ t1). \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ t2) = (P \ t1. \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ t2)$

$\langle \text{proof} \rangle$

lemma *iWeakSince-conj-not*: $\text{finite } I \implies$

$((P \ t1 \wedge \neg Q \ t1). \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ t2) = (P \ t1. \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ t2)$

$\langle \text{proof} \rangle$

lemma *iNextStrong-imp-iNextWeak*: $(\bigcirc_S \ t \ t0 \ I. \ P \ t) \longrightarrow (\bigcirc_W \ t \ t0 \ I. \ P \ t)$

$\langle \text{proof} \rangle$

lemma *iLastStrong-imp-iLastWeak*: $(\ominus_S \ t \ t0 \ I. \ P \ t) \longrightarrow (\ominus_W \ t \ t0 \ I. \ P \ t)$

$\langle \text{proof} \rangle$

lemma *infin-imp-iNextWeak-iNextStrong-eq-iNext*:

$\llbracket \text{infinite } I; \ t0 \in I \rrbracket \implies$

$((\bigcirc_W \ t \ t0 \ I. \ P \ t) = (\bigcirc \ t \ t0 \ I. \ P \ t)) \wedge ((\bigcirc_S \ t \ t0 \ I. \ P \ t) = (\bigcirc \ t \ t0 \ I. \ P \ t))$

$\langle \text{proof} \rangle$

lemma *infin-imp-iNextWeak-eq-iNext*: $\llbracket \text{infinite } I; \ t0 \in I \rrbracket \implies (\bigcirc_W \ t \ t0 \ I. \ P \ t) = (\bigcirc \ t \ t0 \ I. \ P \ t)$

$\langle \text{proof} \rangle$

lemma *infin-imp-iNextStrong-eq-iNext*: $\llbracket \text{infinite } I; \ t0 \in I \rrbracket \implies (\bigcirc_S \ t \ t0 \ I. \ P \ t) = (\bigcirc \ t \ t0 \ I. \ P \ t)$

$\langle \text{proof} \rangle$

lemma *infin-imp-iNextStrong-eq-iNextWeak*: $\llbracket \text{infinite } I; \ t0 \in I \rrbracket \implies (\bigcirc_S \ t \ t0 \ I. \ P \ t) = (\bigcirc_W \ t \ t0 \ I. \ P \ t)$

$\langle \text{proof} \rangle$

lemma

not-in-iNext-eq: $t0 \notin I \implies (\bigcirc \ t \ t0 \ I. \ P \ t) = (P \ t0)$ **and**

not-in-iLast-eq: $t0 \notin I \implies (\ominus \ t \ t0 \ I. \ P \ t) = (P \ t0)$

$\langle \text{proof} \rangle$

lemma

not-in-iNextWeak-eq: $t0 \notin I \implies (\bigcirc_W \ t \ t0 \ I. \ P \ t)$ **and**

not-in-iLastWeak-eq: $t0 \notin I \implies (\ominus_W \ t \ t0 \ I. \ P \ t)$

$\langle \text{proof} \rangle$

lemma

not-in-iNextStrong-eq: $t0 \notin I \implies \neg (\odot_S t t0 I. P t)$ **and**
not-in-iLastStrong-eq: $t0 \notin I \implies \neg (\ominus_S t t0 I. P t)$
 $\langle \text{proof} \rangle$

lemma

iNext-UNIV: $(\odot t t0 UNIV. P t) = P (Suc t0)$ **and**
iNextWeak-UNIV: $(\odot_W t t0 UNIV. P t) = P (Suc t0)$ **and**
iNextStrong-UNIV: $(\odot_S t t0 UNIV. P t) = P (Suc t0)$
 $\langle \text{proof} \rangle$

lemma

iLast-UNIV: $(\ominus t t0 UNIV. P t) = P (t0 - Suc 0)$ **and**
iLastWeak-UNIV: $(\ominus_W t t0 UNIV. P t) = (\text{if } 0 < t0 \text{ then } P (t0 - Suc 0) \text{ else True})$ **and**
iLastStrong-UNIV: $(\ominus_S t t0 UNIV. P t) = (\text{if } 0 < t0 \text{ then } P (t0 - Suc 0) \text{ else False})$
 $\langle \text{proof} \rangle$

lemmas *iTL-Next-UNIV* =

iNext-UNIV iNextWeak-UNIV iNextStrong-UNIV
iLast-UNIV iLastWeak-UNIV iLastStrong-UNIV

lemma *inext-nth-iNext-Suc*: $(\odot t (I \rightarrow n) I. P t) = P (I \rightarrow Suc n)$
 $\langle \text{proof} \rangle$

lemma *iprev-nth-iLast-Suc*: $(\ominus t (I \leftarrow n) I. P t) = P (I \leftarrow Suc n)$
 $\langle \text{proof} \rangle$

3.4 Temporal operators and arithmetic interval operators

Shifting intervals through addition and subtraction of constants. Mirroring intervals through subtraction of intervals from constants. Expanding and compressing intervals through multiplication and division by constants.

Always operator

lemma *iT-Plus-iAll-conv*: $(\Box t I \oplus k. P t) = (\Box t I. P (t + k))$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-iAll-conv*: $(\Box t I \otimes k. P t) = (\Box t I. P (t * k))$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-iAll-conv*: $(\Box t I \oplus - k. P t) = (\Box t (I \downarrow \geq k). P (t - k))$
 $\langle \text{proof} \rangle$

lemma *iT-Minus-iAll-conv*: $(\Box t k \ominus I. P t) = (\Box t (I \downarrow \leq k). P (k - t))$
 $\langle \text{proof} \rangle$

lemma *iT-Div-iAll-conv*: $(\Box t I \odot k. P t) = (\Box t I. P (t \text{ div } k))$
 $\langle \text{proof} \rangle$

lemmas *iT-arith-iAll-conv* =
iT-Plus-iAll-conv
iT-Mult-iAll-conv
iT-Plus-neg-iAll-conv
iT-Minus-iAll-conv
iT-Div-iAll-conv

Eventually operator

lemma
iT-Plus-iEx-conv: $(\Diamond t I \oplus k. P t) = (\Diamond t I. P (t + k))$ **and**
iT-Mult-iEx-conv: $(\Diamond t I \otimes k. P t) = (\Diamond t I. P (t * k))$ **and**
iT-Plus-neg-iEx-conv: $(\Diamond t I \oplus - k. P t) = (\Diamond t (I \downarrow \geq k). P (t - k))$ **and**
iT-Minus-iEx-conv: $(\Diamond t k \ominus I. P t) = (\Diamond t (I \downarrow \leq k). P (k - t))$ **and**
iT-Div-iEx-conv: $(\Diamond t I \odot k. P t) = (\Diamond t I. P (t \text{ div } k))$
 $\langle \text{proof} \rangle$

Until and Since operators

lemma *iT-Plus-iUntil-conv*: $(P t1. t1 \mathcal{U} t2 (I \oplus k). Q t2) = (P (t1 + k). t1 \mathcal{U} t2 I. Q (t2 + k))$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-iUntil-conv*: $(P t1. t1 \mathcal{U} t2 (I \otimes k). Q t2) = (P (t1 * k). t1 \mathcal{U} t2 I. Q (t2 * k))$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-iUntil-conv*: $(P t1. t1 \mathcal{U} t2 (I \oplus - k). Q t2) = (P (t1 - k). t1 \mathcal{U} t2 (I \downarrow \geq k). Q (t2 - k))$
 $\langle \text{proof} \rangle$

lemma *iT-Minus-iUntil-conv*: $(P t1. t1 \mathcal{U} t2 (k \ominus I). Q t2) = (P (k - t1). t1 \mathcal{S} t2 (I \downarrow \leq k). Q (k - t2))$
 $\langle \text{proof} \rangle$

lemma *iT-Div-iUntil-conv*: $(P t1. t1 \mathcal{U} t2 (I \odot k). Q t2) = (P (t1 \text{ div } k). t1 \mathcal{U} t2 I. Q (t2 \text{ div } k))$
 $\langle \text{proof} \rangle$

Until and Since operators can be converted into each other through subtraction of intervals from constants

lemma *iUntil-iSince-conv*:
 $\llbracket \text{finite } I; \text{Max } I \leq k \rrbracket \implies$
 $(P t1. t1 \mathcal{U} t2 I. Q t2) = (P (k - t1). t1 \mathcal{S} t2 (k \ominus I). Q (k - t2))$
 $\langle \text{proof} \rangle$

lemma *iSince-iUntil-conv*:
 $\llbracket \text{finite } I; \text{Max } I \leq k \rrbracket \implies$

$(P \ t1. \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ t2) = (P \ (k - t1). \ t1 \ \mathcal{U} \ t2 \ (k \ominus I). \ Q \ (k - t2))$
 $\langle proof \rangle$

lemma *iT-Plus-iSince-conv*: $(P \ t1. \ t1 \ \mathcal{S} \ t2 \ (I \oplus k). \ Q \ t2) = (P \ (t1 + k). \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ (t2 + k))$
 $\langle proof \rangle$

lemma *iT-Mult-iSince-conv*: $0 < k \implies (P \ t1. \ t1 \ \mathcal{S} \ t2 \ (I \otimes k). \ Q \ t2) = (P \ (t1 * k). \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ (t2 * k))$
 $\langle proof \rangle$

lemma *iT-Plus-neg-iSince-conv*: $(P \ t1. \ t1 \ \mathcal{S} \ t2 \ (I \oplus - k). \ Q \ t2) = (P \ (t1 - k). \ t1 \ \mathcal{S} \ t2 \ (I \downarrow \geq k). \ Q \ (t2 - k))$
 $\langle proof \rangle$

lemma *iT-Minus-iSince-conv*:
 $(P \ t1. \ t1 \ \mathcal{S} \ t2 \ (k \ominus I). \ Q \ t2) = (P \ (k - t1). \ t1 \ \mathcal{U} \ t2 \ (I \downarrow \leq k). \ Q \ (k - t2))$
 $\langle proof \rangle$

lemma *iT-Div-iSince-conv*:
 $0 < k \implies (P \ t1. \ t1 \ \mathcal{S} \ t2 \ (I \oslash k). \ Q \ t2) = (P \ (t1 \text{ div } k). \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ (t2 \text{ div } k))$
 $\langle proof \rangle$

Weak Until and Weak Since operators

lemma *iT-Plus-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (I \oplus k). \ Q \ t2) = (P \ (t1 + k). \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ (t2 + k))$
 $\langle proof \rangle$

lemma *iT-Mult-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (I \otimes k). \ Q \ t2) = (P \ (t1 * k). \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ (t2 * k))$
 $\langle proof \rangle$

lemma *iT-Plus-neg-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (I \oplus - k). \ Q \ t2) = (P \ (t1 - k). \ t1 \ \mathcal{W} \ t2 \ (I \downarrow \geq k). \ Q \ (t2 - k))$
 $\langle proof \rangle$

lemma *iT-Minus-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (k \ominus I). \ Q \ t2) = (P \ (k - t1). \ t1 \ \mathcal{B} \ t2 \ (I \downarrow \leq k). \ Q \ (k - t2))$
 $\langle proof \rangle$

lemma *iT-Div-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (I \oslash k). \ Q \ t2) = (P \ (t1 \text{ div } k). \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ (t2 \text{ div } k))$
 $\langle proof \rangle$

lemma *iT-Plus-iWeakSince-conv*: $(P \ t1. \ t1 \ \mathcal{B} \ t2 \ (I \oplus k). \ Q \ t2) = (P \ (t1 + k). \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ (t2 + k))$
 $\langle proof \rangle$

lemma *iT-Mult-iWeakSince-conv*: $0 < k \implies (P \ t1. \ t1 \ \mathcal{B} \ t2 \ (I \otimes k). \ Q \ t2) = (P \ (t1 * k). \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ (t2 * k))$
 $\langle proof \rangle$

lemma *iT-Plus-neg-iWeakSince-conv*: $(P \ t1. \ t1 \ \mathcal{B} \ t2 \ (I \oplus - k). \ Q \ t2) = (P \ (t1 - k). \ t1 \ \mathcal{B} \ t2 \ (I \downarrow \geq k). \ Q \ (t2 - k))$
 $\langle proof \rangle$

lemma *iT-Minus-iWeakSince-conv*:
 $(P \ t1. \ t1 \ \mathcal{B} \ t2 \ (k \ominus I). \ Q \ t2) = (P \ (k - t1). \ t1 \ \mathcal{W} \ t2 \ (I \downarrow \leq k). \ Q \ (k - t2))$
 $\langle proof \rangle$

lemma *iT-Div-iWeakSince-conv*:
 $0 < k \implies (P \ t1. \ t1 \ \mathcal{B} \ t2 \ (I \oslash k). \ Q \ t2) = (P \ (t1 \ \text{div} \ k). \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ (t2 \ \text{div} \ k))$
 $\langle proof \rangle$

Release and Trigger operators

lemma *iT-Plus-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (I \oplus k). \ Q \ t2) = (P \ (t1 + k). \ t1 \ \mathcal{R} \ t2 \ I. \ Q \ (t2 + k))$
 $\langle proof \rangle$

lemma *iT-Mult-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (I \otimes k). \ Q \ t2) = (P \ (t1 * k). \ t1 \ \mathcal{R} \ t2 \ I. \ Q \ (t2 * k))$
 $\langle proof \rangle$

lemma *iT-Plus-neg-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (I \oplus - k). \ Q \ t2) = (P \ (t1 - k). \ t1 \ \mathcal{R} \ t2 \ (I \downarrow \geq k). \ Q \ (t2 - k))$
 $\langle proof \rangle$

lemma *iT-Minus-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (k \ominus I). \ Q \ t2) = (P \ (k - t1). \ t1 \ \mathcal{T} \ t2 \ (I \downarrow \leq k). \ Q \ (k - t2))$
 $\langle proof \rangle$

lemma *iT-Div-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (I \oslash k). \ Q \ t2) = (P \ (t1 \ \text{div} \ k). \ t1 \ \mathcal{R} \ t2 \ I. \ Q \ (t2 \ \text{div} \ k))$
 $\langle proof \rangle$

lemma *iT-Plus-iTrigger-conv*: $(P \ t1. \ t1 \ \mathcal{T} \ t2 \ (I \oplus k). \ Q \ t2) = (P \ (t1 + k). \ t1 \ \mathcal{T} \ t2 \ I. \ Q \ (t2 + k))$
 $\langle proof \rangle$

lemma *iT-Mult-iTrigger-conv*: $0 < k \implies (P \ t1. \ t1 \ \mathcal{T} \ t2 \ (I \otimes k). \ Q \ t2) = (P \ (t1 * k). \ t1 \ \mathcal{T} \ t2 \ I. \ Q \ (t2 * k))$
 $\langle proof \rangle$

lemma *iT-Plus-neg-iTrigger-conv*: $(P \ t1. \ t1 \ \mathcal{T} \ t2 \ (I \oplus - k). \ Q \ t2) = (P \ (t1 -$

$k). t1 \mathcal{T} t2 (I \downarrow \geq k). Q (t2 - k))$
 $\langle proof \rangle$

lemma *iT-Minus-iTrigger-conv:*

$(P t1. t1 \mathcal{T} t2 (k \ominus I). Q t2) = (P (k - t1). t1 \mathcal{R} t2 (I \downarrow \leq k). Q (k - t2))$
 $\langle proof \rangle$

lemma *iT-Div-iTrigger-conv:*

$0 < k \implies (P t1. t1 \mathcal{T} t2 (I \oslash k). Q t2) = (P (t1 \text{ div } k). t1 \mathcal{T} t2 I. Q (t2 \text{ div } k))$
 $\langle proof \rangle$

end