

Morley's Theorem*

Benjamin Puyobro

Université Paris Saclay, IRT SystemX, LMF, CNRS

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theory *Complex-Angles*

imports *Complex-Geometry.Elementary-Complex-Geometry*

begin

1 Introduction

Morley's trisector theorem states that in any triangle, the three points of intersection of the adjacent angle trisectors form an equilateral triangle, called the first Morley triangle or simply the Morley triangle. This theorem is listed in the 100 theorem list [3].

In this theory, we define some basics elements on complex geometry such as axial symmetry, rotations. We also define basic property of congruent triangles in the complex field following the model already presented in the afp [1] In addition we demonstrate sines law in the complex context.

Finally we prove the Morley theorem using Alain Connes's proof [2].

2 Angles between three complex

definition *angle-c-def*: $\langle \text{angle-c } a \ b \ c \equiv \angle (a-b) (c-b) \rangle$

lemma *angle-c-commute-pi*:

assumes $\langle \text{angle-c } a \ b \ c = \text{pi} \rangle$

shows $\text{angle-c } a \ b \ c = \text{angle-c } c \ b \ a$

using *angle-c-def* *assms* *ang-vec-sym-pi* **by** *presburger*

lemma *angle-c-commute*:

assumes $\langle \text{angle-c } a \ b \ c \neq \text{pi} \rangle$

shows $\text{angle-c } a \ b \ c = -\text{angle-c } c \ b \ a$

using *angle-c-def* *assms* *ang-vec-sym* **by** *presburger*

lemma *angle-c-sum*:

assumes $1: \langle a - b \neq 0 \rangle$ **and** $2: \langle c - b \neq 0 \rangle$ **and** $3: \langle b - d \neq 0 \rangle$

shows $\langle \text{angle-c } a \ b \ c + \text{angle-c } c \ b \ d \rangle = \text{angle-c } a \ b \ d$

using *angle-c-def* *canon-ang-sum* **by** *force*

lemma *collinear-angle*: $\langle \text{collinear } a \ b \ c \implies a \neq b \implies b \neq c \implies c \neq a \implies \text{angle-c } a \ b \ c = 0 \vee \text{angle-c } a \ b \ c = \text{pi} \rangle$

unfolding *collinear-def* **unfolding** *angle-c-def*

by (*metis* *ang-vec-def* *arg-div* *collinear-def* *collinear-sym2'* *is-real-arg2* *right-minus-eq*)

lemma *cdist-commute*: $\langle \text{cdist } a \ b = \text{cdist } b \ a \rangle$

by (*simp* *add*: *norm-minus-commute*)

lemma *angle-sum-triangle*:
assumes $h: a \neq b \wedge b \neq c \wedge a \neq c$
shows $\lfloor \angle (c-a) (b-a) + \angle (a-b) (c-b) + \angle (b-c) (a-c) \rfloor = \pi$
proof –
have $f0: \langle b-a \neq 0 \rangle \langle c-a \neq 0 \rangle \langle b-c \neq 0 \rangle$ **using** h **by** *auto*
then have $\langle \angle c (c-a) (b-a) = \text{abs} (\text{Arg} ((b-a)/(c-a))) \rangle$
unfolding *ang-vec-def ang-vec-c-def*
apply (*subst arg-div* [*OF f0(1) f0(2), symmetric*])
by (*simp*)
have $\langle (b-a)/(c-a) = \text{rcis} (\text{cmod} ((b-a)/(c-a))) (\text{Arg} ((b-a)/(c-a))) \rangle$
by (*simp add: rcis-cmod-Arg*)
have $\langle (b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c) = -1 \rangle$
using h
by (*smt (z3) divide-eq-0-iff divide-eq-minus-1-iff f0(1,2,3) minus-diff-eq mult-minus-right nonzero-eq-divide-eq times-divide-eq-left*)
then have $\langle \text{Arg} ((b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c)) = \pi \rangle$
by (*metis arg-cis cis-pi pi-canonical*)
then have $\langle \text{Arg} ((b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c))$
 $= \lfloor \text{Arg} ((b-a)/(c-a) * (c-b)/(a-b)) + \text{Arg} ((a-c)/(b-c)) \rfloor \rangle$
using *arg-mult f0*
by (*metis (no-types, lifting) Arg-zero arg-mult divide-divide-eq-left' pi-neq-zero times-divide-eq-left times-divide-eq-right*)
then have $\langle \lfloor \text{Arg} ((b-a)/(c-a) * (c-b)/(a-b)) + \text{Arg} ((a-c)/(b-c)) \rfloor$
 $= \lfloor \text{Arg} ((b-a)/(c-a)) + \text{Arg} ((c-b)/(a-b)) + \text{Arg} ((a-c)/(b-c)) \rfloor \rangle$
by (*metis (no-types, lifting) arg-mult canon-ang-arg canon-ang-sum divide-eq-0-iff eq-iff-diff-eq-0 h no-zero-divisors times-divide-eq-right*)
then show *?thesis*
using
 $\langle \text{Arg} ((b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c)) = \lfloor \text{Arg} ((b-a)/(c-a) * (c-b)/(a-b)) + \text{Arg} ((a-c)/(b-c)) \rfloor \rangle$
 $\langle \text{Arg} ((b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c)) = \pi \rangle$ *arg-div*
 h **by** *force*
qed

lemma *angle-sum-triangle-c*:
assumes $h: a \neq b \wedge b \neq c \wedge a \neq c$
shows $\lfloor \text{angle-c } c \ a \ b + \text{angle-c } a \ b \ c + \text{angle-c } b \ c \ a \rfloor = \pi$
unfolding *angle-c-def*
using h **by** (*smt (z3) angle-sum-triangle h*)

lemma *angle-sum-triangle'*:
assumes $h: a \neq 0 \wedge b \neq 0 \wedge c \neq 0$
shows $\lfloor \angle (-a) \ b + \angle (-b) \ c + \angle (-c) \ a \rfloor = \pi$
proof –
have $f0: \langle a \neq 0 \rangle \langle b \neq 0 \rangle \langle c \neq 0 \rangle$ **using** h **by** *auto*
have $\langle b/(-a) * a/(-c) * c/(-b) = -1 \rangle$

```

    using h by(auto)
    then have f10:⟨Arg (b/(-a) * a/(-c) * c/(-b)) = pi⟩
      by (metis arg-cis cis-pi pi-canonical)
    then have f1:⟨Arg (b/(-a) * a/(-c) * c/(-b))
      = ⟦Arg (b/(-a) * a/(-c)) + Arg (c/(-b))⟧⟩
      using arg-mult f0
    by (metis ⟨b / - a * a / - c * c / - b = - 1⟩ divide-eq-0-iff divide-eq-minus-1-iff
      times-divide-eq-right)
    then have f2:⟨⟦Arg (b/(-a) * a/(-c)) + Arg (c/(-b))⟧
      = ⟦Arg (b/(-a)) + Arg (a/(-c)) + Arg (c/(-b))⟧⟩
      by (metis (no-types, lifting) arg-mult canon-ang-arg canon-ang-sum divide-eq-0-iff
        f0(1,2,3)
        neg-0-equal-iff-equal no-zero-divisors times-divide-eq-right)
    then show ?thesis
      using f1 f2 arg-div h
      by (smt (verit) f10 add.left-commute ang-vec-def neg-equal-0-iff-equal)
qed

lemma ang-c-in:⟨angle-c a b c ∈ {-pi<..

```

3 Complex triangles

In this section we define triangles and derive some useful lemmas on congruent triangles, following the model [1]

```

locale congruent-ctriangle =
  fixes a1 b1 c1 :: complex and a2 b2 c2 :: complex
  assumes sides': cdist a1 b1 = cdist a2 b2 cdist a1 c1 = cdist a2 c2 cdist b1 c1
    = cdist b2 c2
  and angles': angle-c b1 a1 c1 = angle-c b2 a2 c2 ∨ angle-c b1 a1 c1 = -
    angle-c b2 a2 c2

```

$\text{angle-c } a1 \ b1 \ c1 = \text{angle-c } a2 \ b2 \ c2 \vee \text{angle-c } a1 \ b1 \ c1 = - \text{angle-c } a2 \ b2 \ c2$
 $\text{angle-c } a1 \ c1 \ b1 = \text{angle-c } a2 \ c2 \ b2 \vee \text{angle-c } a1 \ c1 \ b1 = - \text{angle-c } a2 \ c2 \ b2$
begin

lemma *sides*:

$\text{cdist } a1 \ b1 = \text{cdist } a2 \ b2 \text{ cdist } a1 \ c1 = \text{cdist } a2 \ c2 \text{ cdist } b1 \ c1 = \text{cdist } b2 \ c2$
 $\text{cdist } b1 \ a1 = \text{cdist } a2 \ b2 \text{ cdist } c1 \ a1 = \text{cdist } a2 \ c2 \text{ cdist } c1 \ b1 = \text{cdist } b2 \ c2$
 $\text{cdist } a1 \ b1 = \text{cdist } b2 \ a2 \text{ cdist } a1 \ c1 = \text{cdist } c2 \ a2 \text{ cdist } b1 \ c1 = \text{cdist } c2 \ b2$
 $\text{cdist } b1 \ a1 = \text{cdist } b2 \ a2 \text{ cdist } c1 \ a1 = \text{cdist } c2 \ a2 \text{ cdist } c1 \ b1 = \text{cdist } c2 \ b2$
using *sides'*
by (*auto simp add: norm-minus-commute*)

lemma *angles*:

$\text{angle-c } b1 \ a1 \ c1 = \text{angle-c } b2 \ a2 \ c2 \vee \text{angle-c } b1 \ a1 \ c1 = - \text{angle-c } b2 \ a2 \ c2$
 $\text{angle-c } a1 \ b1 \ c1 = \text{angle-c } a2 \ b2 \ c2 \vee \text{angle-c } a1 \ b1 \ c1 = - \text{angle-c } a2 \ b2 \ c2$
 $\text{angle-c } a1 \ c1 \ b1 = \text{angle-c } a2 \ c2 \ b2 \vee \text{angle-c } a1 \ c1 \ b1 = - \text{angle-c } a2 \ c2 \ b2$
 $\text{angle-c } c1 \ a1 \ b1 = \text{angle-c } b2 \ a2 \ c2 \vee \text{angle-c } c1 \ a1 \ b1 = - \text{angle-c } b2 \ a2 \ c2$
 $\text{angle-c } c1 \ b1 \ a1 = \text{angle-c } a2 \ b2 \ c2 \vee \text{angle-c } c1 \ b1 \ a1 = - \text{angle-c } a2 \ b2 \ c2$
 $\text{angle-c } b1 \ c1 \ a1 = \text{angle-c } a2 \ c2 \ b2 \vee \text{angle-c } b1 \ c1 \ a1 = - \text{angle-c } a2 \ c2 \ b2$
 $\text{angle-c } b1 \ a1 \ c1 = \text{angle-c } c2 \ a2 \ b2 \vee \text{angle-c } b1 \ a1 \ c1 = - \text{angle-c } c2 \ a2 \ b2$
 $\text{angle-c } a1 \ b1 \ c1 = \text{angle-c } c2 \ b2 \ a2 \vee \text{angle-c } a1 \ b1 \ c1 = - \text{angle-c } c2 \ b2 \ a2$
 $\text{angle-c } a1 \ c1 \ b1 = \text{angle-c } b2 \ c2 \ a2 \vee \text{angle-c } a1 \ c1 \ b1 = - \text{angle-c } b2 \ c2 \ a2$
 $\text{angle-c } c1 \ a1 \ b1 = \text{angle-c } c2 \ a2 \ b2 \vee \text{angle-c } c1 \ a1 \ b1 = - \text{angle-c } c2 \ a2 \ b2$
 $\text{angle-c } c1 \ b1 \ a1 = \text{angle-c } c2 \ b2 \ a2 \vee \text{angle-c } c1 \ b1 \ a1 = - \text{angle-c } c2 \ b2 \ a2$
 $\text{angle-c } b1 \ c1 \ a1 = \text{angle-c } b2 \ c2 \ a2 \vee \text{angle-c } b1 \ c1 \ a1 = - \text{angle-c } b2 \ c2 \ a2$
using *angles' angle-c-commute*
by(*simp-all add: angle-c-commute*)(*metis add.inverse-inverse angle-c-commute*)+

end

end

theory *Complex-Trigonometry*

imports *HOL-Analysis.Convex Complex-Angles Complex-Triangles-Definitions*

begin

4 Complex trigonometry

In this section we add some trigonometric results, especially the law of sines

lemma *ang-sin*:

shows $\langle \text{Im } ((b-a) * \text{cnj}(c-a)) = \text{cmod } (c-a) * \text{cmod } (b-a) * \sin (\text{angle-c } c \ a \ b) \rangle$
proof –
have $f0: \langle \sin (\text{Arg } (\text{cnj } c - \text{cnj } a)) = - \sin (\text{Arg } (c - a)) \rangle$
by (*metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff neg-0-equal-iff-equal sin-minus sin-pi*)

have $f1: \langle \cos (\text{Arg } (\text{cnj } c - \text{cnj } a)) = \cos (\text{Arg } (c - a)) \rangle$
by (*metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff cos-minus*)
then have $\langle b-a = \text{cmod } (b-a) * \text{cis } (\text{Arg } (b-a)) \wedge \text{cnj } (c-a) = \text{cmod } (c-a) * \text{cis } (\text{Arg } (\text{cnj } (c-a))) \rangle$
using *rcis-cmod-Arg rcis-def*
by (*metis complex-mod-cnj*)
then have $\langle \text{Im}(\text{cis } (\text{Arg } (b-a)) * \text{cis } (\text{Arg } (\text{cnj } (c-a)))) = \sin (\text{angle-c } c \ a \ b) \rangle$
unfolding *ang-vec-def angle-c-def* **using** $f1 \ f0$ **by** (*auto simp: cis.code sin-diff*)
then have $\langle \text{Im} (\text{cmod } (b-a) * \text{cis } (\text{Arg } (b-a)) * \text{cmod } (c-a) * \text{cis } (\text{Arg } (\text{cnj } (c-a)))) = \text{cmod } (c-a) * \text{cmod } (b-a) * \sin (\text{angle-c } c \ a \ b) \rangle$
by (*metis Im-complex-of-real[of cmod (c - a)] Im-complex-of-real[of cmod (b - a)]*)
 $\text{Im-mult-real[of cor (cmod (c - a)) \text{cis } (\text{Arg } (\text{cnj } (c - a))) * (\text{cor (cmod } (b - a)) * \text{cis } (\text{Arg } (b - a)))]$
 $\text{Im-mult-real[of cor (cmod (b - a)) \text{cis } (\text{Arg } (b - a)) * \text{cis } (\text{Arg } (\text{cnj } (c - a)))]$
 $\text{Re-complex-of-real[of cmod (c - a)] \text{Re-complex-of-real[of cmod (b - a)]}$
 $\text{ab-semigroup-mult-class.mult-ac}(1)[\text{of cor (cmod (b - a)) \text{cis } (\text{Arg } (b - a)) \text{cis } (\text{Arg } (\text{cnj } (c - a)))]$
 $\text{ab-semigroup-mult-class.mult-ac}(1)[\text{of cis } (\text{Arg } (\text{cnj } (c - a))) \text{cor (cmod (b - a)) * cis } (\text{Arg } (b - a)) \text{cor (cmod (c - a))}]$
 $\text{ab-semigroup-mult-class.mult-ac}(1)[\text{of cmod (c - a) cmod (b - a) Im (cis (Arg (b - a)) * cis (Arg (cnj (c - a))))}]$
 $\text{mult.commute[of cis (Arg (cnj (c - a))) \text{cor (cmod (b - a)) * cis (Arg (b - a))}]$
 $\text{mult.commute[of cis (Arg (cnj (c - a))) * (\text{cor (cmod (b - a)) * cis (Arg (b - a))}]$
 $\text{cor (cmod (c - a))}]$
 $\text{mult.commute[of cor (cmod (b - a)) * cis (Arg (b - a)) * cor (cmod (c - a)) \text{cis (Arg (cnj (c - a)))]}$
then show $\langle \text{Im } ((b-a) * \text{cnj } (c-a)) = \text{cmod } (c-a) * \text{cmod } (b-a) * \sin (\text{angle-c } c \ a \ b) \rangle$
by (*metis* $\langle b-a = \text{cor } (\text{cmod } (b-a)) * \text{cis } (\text{Arg } (b-a)) \wedge \text{cnj } (c-a) = \text{cor } (\text{cmod } (c-a)) * \text{cis } (\text{Arg } (\text{cnj } (c-a))) \rangle$)
ab-semigroup-mult-class.mult-ac(1))
qed

lemma *ang-cos*:

shows $\langle \text{Re } ((b-a) * \text{cnj } (c-a)) = \text{cmod } (c-a) * \text{cmod } (b-a) * \cos (\text{angle-c } c \ a \ b) \rangle$
proof –
have $f0: \langle \sin (\text{Arg } (\text{cnj } c - \text{cnj } a)) = - \sin (\text{Arg } (c - a)) \rangle$
by (*metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff neg-0-equal-iff-equal sin-minus sin-pi*)
have $f1: \langle \cos (\text{Arg } (\text{cnj } c - \text{cnj } a)) = \cos (\text{Arg } (c - a)) \rangle$
by (*metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff cos-minus*)
then have $f2: \langle b-a = \text{cmod } (b-a) * \text{cis } (\text{Arg } (b-a)) \wedge \text{cnj } (c-a) = \text{cmod } (c-a) * \text{cis } (\text{Arg } (\text{cnj } (c-a))) \rangle$

```

    using rcis-cmod-Arg rcis-def
    by (metis complex-mod-cnj)
    then have ⟨Re(cis (Arg (b-a)) * cis (Arg (cnj(c-a)))) = cos (angle-c c a b)⟩
    unfolding ang-vec-def angle-c-def using f1 f0 by(auto simp:cis.code cos-diff)
    then have ⟨Re (cmod (b-a) * cmod (c-a)* cis (Arg (b-a)) * cis (Arg (cnj
(c-a)))) =
      Re (cmod (c-a) *cmod (b-a)) * cos (angle-c c a b) ⟩
    proof -
      have f1: cor (cmod (b - a) * cmod (c - a)) * cis (Arg (b - a)) * cis (Arg
(cnj (c - a)))
    = cor (cmod (b - a) * cmod (c - a)) * (cis (Arg (b - a)) * cis (Arg (cnj (c -
a))))
      using ab-semigroup-mult-class.mult-ac(1) by blast
      have cmod (b - a) * cmod (c - a) = cmod (c - a) * cmod (b - a)
      by argo
      then show ?thesis
      using f1 Im-complex-of-real Re-mult-real ⟨Re (cis (Arg (b - a)) * cis (Arg
(cnj (c - a))))⟩
      = cos (angle-c c a b)⟩ by presburger
    qed
    then show ⟨Re ((b-a)*cnj(c-a)) = cmod (c-a) *cmod (b-a) * cos (angle-c c
a b)⟩
    by (metis Re-complex-of-real f2 mult.assoc mult.commute of-real-mult)

```

qed

lemma *law-of-cosines'*:

```

    assumes h:⟨A≠C⟩ ⟨A≠B⟩
    shows ((cdist B C)2 - (cdist A C)2 - (cdist A B)2) / (- 2*(cdist A C)*(cdist
A B)) = (cos (∠ (C-A) (B-A)))
    using law-of-cosines[of B C A] h by(auto simp:field-simps)

```

lemma *law-of-cosines''*:

```

    shows (cdist A C)2 = (cdist B C)2 - (cdist A B)2 + 2*(cdist A C)*(cdist A
B)*(cos (∠ (C-A) (B-A)))
    using law-of-cosines[of B C A] by(auto simp:field-simps)

```

lemma *law-of-cosines'''*:

```

    shows (cdist A B)2 = (cdist B C)2 - (cdist A C)2 + 2*(cdist A C)*(cdist A
B)*(cos (∠ (C-A) (B-A)))
    using law-of-cosines[of B C A] by(auto simp:field-simps)

```

4.1 The law of sines

theorem *law-of-sines*:

```

    assumes h1:⟨b ≠ a⟩ ⟨a ≠ c⟩ ⟨b ≠ c⟩
    shows sin (angle-c a b c) * cdist b c = sin (angle-c c a b) * cdist a c (is ?A =
?B)
    proof -

```

```

{fix a b c :: complex
  have f0:⟨sin (Arg (cnj c - cnj a)) = - sin (Arg (c - a))⟩
  by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff neg-0-equal-iff-equal sin-minus
      sin-pi)
  have f1:⟨cos (Arg (cnj c - cnj a)) = cos (Arg (c - a))⟩
  by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff cos-minus)
  assume ⟨b-a ≠ 0⟩ ⟨c-a ≠ 0⟩
  then have f2: ⟨b-a = cmod (b-a) * cis (Arg (b-a)) ∧ cnj (c-a) = cmod
(c-a) * cis (Arg (cnj (c-a)))⟩
    using rcis-cmod-Arg rcis-def
    by (metis complex-mod-cnj)
  then have ⟨Im (cis (Arg (b-a)) * cis (Arg (cnj (c-a)))) = sin (angle-c c a b)⟩
    unfolding ang-vec-def angle-c-def using f1 f0 by (auto simp: cis.code sin-diff)

  then have ⟨Im (cmod (b-a) * cis (Arg (b-a)) * cmod (c-a) * cis (Arg (cnj
(c-a)))) =
    cmod (c-a) * cmod (b-a) * sin (angle-c c a b)⟩
    by (metis Im-complex-of-real[of cmod (c - a)] Im-complex-of-real[of cmod (b
- a)]
      Im-mult-real[of cor (cmod (c - a))
        cis (Arg (cnj (c - a))) * (cor (cmod (b - a)) * cis (Arg (b - a)))]
      Im-mult-real[of cor (cmod (b - a)) cis (Arg (b - a)) * cis (Arg (cnj (c -
a)))]
      Re-complex-of-real[of cmod (c - a)] Re-complex-of-real[of cmod (b - a)]
      ab-semigroup-mult-class.mult-ac(1)[of cor (cmod (b - a)) cis (Arg (b -
a))
        cis (Arg (cnj (c - a)))]
      ab-semigroup-mult-class.mult-ac(1)[of cis (Arg (cnj (c - a)))]
      cor (cmod (b - a)) * cis (Arg (b - a)) cor (cmod (c - a))
      ab-semigroup-mult-class.mult-ac(1)[of cmod (c - a) cmod (b - a)
        Im (cis (Arg (b - a)) * cis (Arg (cnj (c - a))))]
      mult.commute[of cis (Arg (cnj (c - a))) cor (cmod (b - a)) * cis (Arg (b
- a))
        mult.commute[of cis (Arg (cnj (c - a))) * (cor (cmod (b - a)) * cis (Arg
(b - a))
          cor (cmod (c - a))
          mult.commute[of cor (cmod (b - a)) * cis (Arg (b - a)) * cor (cmod (c
- a))
            cis (Arg (cnj (c - a)))]
        then have ⟨Im ((b-a)*cnj(c-a)) = cmod (c-a) * cmod (b-a) * sin (angle-c
c a b)⟩
        using ang-sin by presburger
      }note ang=this
      have i2:⟨sin (angle-c c a b) = Im((b-a)*cnj(c-a)) / (cmod(b-a)*cmod(c-a))⟩
        using ang[of b a c] h1 by (auto)
      have ⟨sin (angle-c a b c) = Im((c-b)*cnj(a-b)) / (cmod (c-b)*cmod (a-b))⟩
        using ang h1 by (auto)
      then have imp1:⟨cmod (c-b) * sin (angle-c a b c) = Im((c-b)*cnj(a-b)) /
(cmod (a-b))⟩

```



```

    by auto
    from i2 have imp2:⟨cmod (c-a) * sin (angle-c c a b) = Im((b-a)*cnj(c-a))/
(cmod(b-a))⟩
    by auto
    show ?thesis using imp1 imp2 by(auto simp:field-simps norm-minus-commute)

qed

lemma law-of-sines': assumes h1:⟨b ≠ a⟩ ⟨a ≠ c⟩ ⟨b ≠ c⟩
  shows sin (angle-c a b c) * cdist b a = sin (angle-c b c a) * cdist a c
proof -
  {fix a b c ::complex
   have f0:⟨sin (Arg (cnj c - cnj a)) = - sin (Arg (c - a))⟩
   by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff neg-0-equal-iff-equal sin-minus
sin-pi)
   have f1:⟨cos (Arg (cnj c - cnj a)) = cos (Arg (c - a))⟩
   by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff cos-minus)
   assume ⟨b-a ≠ 0⟩ ⟨c-a ≠ 0⟩
   then have f2:⟨b-a = cmod (b-a) * cis (Arg (b-a)) ∧ cnj (c-a) = cmod
(c-a) * cis (Arg (cnj (c-a)))⟩
   using rcis-cmod-Arg rcis-def
   by (metis complex-mod-cnj)
   then have ⟨Im (cis (Arg (b-a)) * cis (Arg (cnj(c-a)))) = sin (angle-c c a b)⟩
   unfolding ang-vec-def angle-c-def using f0 f1 by(auto simp:cis.code sin-diff)

   then have ⟨Im (cmod (b-a) * cis (Arg (b-a)) * cmod (c-a) * cis (Arg (cnj
(c-a)))) =
cmod (c-a) * cmod (b-a) * sin (angle-c c a b)⟩
   by (metis Im-complex-of-real[of cmod (c - a)] Im-complex-of-real[of cmod (b
- a)]
Im-mult-real[of cor (cmod (c - a))
cis (Arg (cnj (c - a))) * (cor (cmod (b - a)) * cis (Arg (b - a)))]
Im-mult-real[of cor (cmod (b - a)) cis (Arg (b - a)) * cis (Arg (cnj (c -
a)))]
Re-complex-of-real[of cmod (c - a)] Re-complex-of-real[of cmod (b - a)]
ab-semigroup-mult-class.mult-ac(1)[of cor (cmod (b - a)) cis (Arg (b -
a))
cis (Arg (cnj (c - a)))]
ab-semigroup-mult-class.mult-ac(1)[of cis (Arg (cnj (c - a)))]
cor (cmod (b - a)) * cis (Arg (b - a)) cor (cmod (c - a))]
ab-semigroup-mult-class.mult-ac(1)[of cmod (c - a) cmod (b - a)
Im (cis (Arg (b - a)) * cis (Arg (cnj (c - a))))]
mult.commute[of cis (Arg (cnj (c - a))) cor (cmod (b - a)) * cis (Arg (b
- a))]
mult.commute[of cis (Arg (cnj (c - a))) * (cor (cmod (b - a)) * cis (Arg
(b - a)))]
cor (cmod (c - a))]
mult.commute[of cor (cmod (b - a)) * cis (Arg (b - a)) * cor (cmod (c
- a))]
```

```

      cis (Arg (cnj (c - a))))]
    then have ⟨Im ((b-a)*cnj(c-a)) = cmod (c-a) * cmod (b-a) * sin (angle-c
c a b)⟩
      using ang-sin by presburger
    }note ang=this
    have i2:⟨sin (angle-c b c a) = Im((a-c)*cnj(b - c)) / (cmod(b-c)*cmod(a-c))⟩
      using ang[of a c b] h1 by(auto)
    have ⟨sin (angle-c a b c) = Im((c-b)*cnj(a-b)) / (cmod (c-b)*cmod (a-b))⟩
      using ang h1 by(auto)
    then have imp1:cmod (a-b) * sin (angle-c a b c) = Im((c-b)*cnj(a-b)) /
(cmod (c-b)) ›
      by auto
    from i2 have imp2:cmod (a-c) * sin (angle-c b c a) = Im((a-c)*cnj(b-c)) /
(cmod(b-c)) ›
      by auto
    show ?thesis using imp1 imp2
      by (metis cdist-commute h1(1,2,3) law-of-sines)
qed

```

```

lemma ang-pos-pos:⟨q≠p ⟹ p≠ r ⟹ r ≠ q ⟹ angle-c q r p ≥ 0 ⟹ angle-c
r p q ≥ 0⟩
  using law-of-sines'[of r q p]
  by (smt (verit) ang-vec-bounded angle-c-def cdist-def mult-neg-pos mult-nonneg-nonneg
right-minus-eq sin-ge-zero sin-gt-zero sin-minus zero-less-norm-iff)

```

```

lemma cmod-pos:⟨cmod a ≥ 0⟩
  by simp

```

```

lemma ang-neg-neg:⟨q≠p ⟹ p≠ r ⟹ r ≠ q ⟹ angle-c q r p < 0 ⟹ angle-c
r p q < 0⟩
proof -
  assume ⟨q≠p ⟩ ⟨p≠ r ⟩ ⟨r ≠ q ⟩ ⟨angle-c q r p < 0⟩
  then have ⟨sin (angle-c q r p) < 0⟩
    using ang-c-in
    by (metis ang-vec-def angle-c-def canon-ang(1) minus-less-iff neg-0-less-iff-less
sin-gt-zero sin-minus)
  then have ⟨sin (angle-c q r p) * cdist r q < 0⟩
    using ⟨r ≠ q ⟩ mult-neg-pos by fastforce
  from law-of-sines'[of r q p] have ⟨sin (angle-c r p q) < 0⟩
    by (metis ⟨p ≠ r ⟩ ⟨q ≠ p ⟩ ⟨r ≠ q ⟩ ⟨sin (angle-c q r p) * cdist r q < 0⟩
cdist-def linorder-not-less mult-less-0-iff norm-ge-zero)
  then show ?thesis
    using ang-c-in[of r p q]
    by (metis ang-vec-def angle-c-def canon-ang(2) linorder-not-less sin-ge-zero)
qed

```

```

lemma collinear-sin-neq-0:

```

$\langle \neg \text{collinear } a2 \ b2 \ c2 \implies \sin (\text{angle-c } a2 \ c2 \ b2) \neq 0 \rangle$
unfolding *collinear-def angle-c-def*
by (*metis Im-complex-div-eq-0 ang-sin angle-c-def collinear-def collinear-sym1 collinear-sym2' mult-eq-0-iff*)

lemma *collinear-sin-neq-pi*:
 $\langle \neg \text{collinear } a2 \ b2 \ c2 \implies \sin (\text{angle-c } a2 \ c2 \ b2) \neq \pi \rangle$
unfolding *collinear-def angle-c-def*
by (*metis add-cancel-right-left order.antisym order.trans le-add-same-cancel1 zero-le-one one-add-one one-neq-zero pi-ge-two sin-le-one*)

lemma *collinear-iff*:
assumes $\langle a \neq b \wedge b \neq c \wedge c \neq a \rangle$
shows $\langle \text{collinear } a \ b \ c \longleftrightarrow (\text{angle-c } a \ b \ c = \pi \vee \text{angle-c } a \ b \ c = 0) \rangle$
using *assms collinear-angle collinear-sin-neq-0 collinear-sym1 sin-minus-pi sin-pi-minus*
by *fastforce*

definition $\langle \text{innerprod } a \ b \equiv \text{cnj } a * b \rangle$

lemma *left-lin-innerprod*: $\langle \text{innerprod } (x + y) \ z = \text{innerprod } x \ z + \text{innerprod } y \ z \rangle$
unfolding *innerprod-def*
by (*simp add: mult.commute ring-class.ring-distrib(1)*)

lemma *right-lin-innerprod*: $\langle \text{innerprod } x \ (y + z) = \text{innerprod } x \ y + \text{innerprod } x \ z \rangle$
unfolding *innerprod-def*
by (*simp add: ring-class.ring-distrib(1)*)

lemma *leftlin-innerprod*: $\langle \text{innerprod } x \ (t * y) = t * \text{innerprod } x \ y \rangle$
unfolding *innerprod-def* **by** (*auto*)

lemma *rightsqliin-innerprod*: $\langle \text{innerprod } (t * x) \ (y) = \text{cnj } t * \text{innerprod } x \ y \rangle$
unfolding *innerprod-def* **by** (*auto*)

lemma *norm-eq-csqrt-inner*: $\langle \text{norm } x = \text{csqrt } (\text{innerprod } x \ x) \rangle$
using *complex-mod-sqrt-Re-mult-cnj innerprod-def* **by** *force*

lemma *abs2-eq-inner*: $\langle \text{abs } (\text{innerprod } x \ y)^2 = \text{innerprod } x \ y * \text{cnj } (\text{innerprod } x \ y) \rangle$
unfolding *innerprod-def abs-complex-def* **apply** (*rule complex-eqI*)
by (*metis comp-apply complex-mult-cnj-cmod of-real-power*)**+**

lemma *complex-add-inner-cnj*: $\langle t * \text{innerprod } x \ y + \text{cnj } (t * \text{innerprod } x \ y) = 2 * \text{Re } (t * \text{innerprod } x \ y) \rangle$
using *complex-add-cnj* **by** *blast*

```

lemma Re-innerprod-inner: $\langle \text{Re } (\text{innerprod } (a-b) (c-b)) = (a-b) \cdot (c-b) \rangle$ 
  unfolding innerprod-def inner-complex-def by (auto simp:field-simps)

lemma angle-c-arccos-pos:
  assumes  $h: \langle a \neq b \wedge b \neq c \rangle \langle \text{angle-c } a \ b \ c \geq 0 \rangle$ 
  shows  $\langle \text{angle-c } a \ b \ c = \arccos ((\text{Re } (\text{innerprod } (a-b) (c-b))) / (\text{cmod}(a-b) * \text{cmod}(c-b))) \rangle$ 
proof -
  have  $\langle \text{Re } ((a-b) * \text{cnj}(c-b)) = (\text{Re } (\text{innerprod } (a-b) (c-b))) \rangle$ 
  unfolding innerprod-def by (auto simp:field-simps)
  then have  $\langle ((\text{Re } (\text{innerprod } (a-b) (c-b))) / (\text{cmod}(a-b) * \text{cmod}(c-b))) = \cos$ 
    (angle-c a b c)  $\rangle$ 
  by (metis (no-types, lifting) ang-cos angle-c-commute cos-minus divisors-zero
    h(1) mult.commute nonzero-mult-div-cancel-left norm-eq-zero right-minus-eq)
  then show ?thesis
  using ang-vec-bounded angle-c-def arccos-cos h(2) by presburger
qed

lemma angle-c-arccos-neg:
  assumes  $h: \langle a \neq b \wedge b \neq c \rangle \langle \text{angle-c } a \ b \ c \leq 0 \rangle$ 
  shows  $\langle - \text{angle-c } a \ b \ c = \arccos ((\text{Re } (\text{innerprod } (a-b) (c-b))) / (\text{cmod}(a-b) * \text{cmod}(c-b))) \rangle$ 
proof -
  have  $\langle \text{Re } ((a-b) * \text{cnj}(c-b)) = (\text{Re } (\text{innerprod } (a-b) (c-b))) \rangle$ 
  unfolding innerprod-def by (auto simp:field-simps)
  then have  $\langle ((\text{Re } (\text{innerprod } (a-b) (c-b))) / (\text{cmod}(a-b) * \text{cmod}(c-b))) = \cos$ 
    (angle-c a b c)  $\rangle$ 
  by (metis (no-types, lifting) ang-cos angle-c-commute cos-minus divisors-zero
    h(1) mult.commute nonzero-mult-div-cancel-left norm-eq-zero right-minus-eq)
  then show ?thesis
  using ang-vec-bounded angle-c-def arccos-cos h(2)
  by (metis arccos-cos2 less-le-not-le)
qed

end
theory Third-Unity-Root

```

imports *Complex-Angles*

begin

5 Third unity root

In this section we prove some basic properties of the third unity root j .

```

lemma root-unity-3:  $\langle (z::\text{complex})^3 - 1 = 0 \longleftrightarrow (z = \text{cis } (2 * \pi / 3) \vee z = 1 \vee z = \text{cis } (4 * \pi / 3)) \rangle$ 
proof (rule iffI)
  have  $\langle \text{cis } (2 * \pi * \text{real } xa / 3) \neq 1 \implies$ 

```

```

      cis (2 * pi * real xa / 3) ≠ cis (2 * pi / 3)
    => xa < 3 => cis (2 * pi * real xa / 3) = cis (4 * pi / 3) ›
  for xa
proof(rule ccontr)
  assume h:⟨cis (2 * pi * real xa / 3) ≠ 1⟩
  ⟨cis (2 * pi * real xa / 3) ≠ cis (2 * pi / 3)⟩
  ⟨xa < 3⟩ ⟨cis (2 * pi * real xa / 3) ≠ cis (4 * pi / 3)⟩
  then have ⟨xa = 1 ∨ xa = 0⟩
  using less-Suc-eq numeral-3-eq-3 by fastforce
  then show False
  using h(1) h(2) by auto
qed
  then have f0:⟨(λk. cis (2 * pi * real k / real 3)) ‘ {.. $3$ } = {1, cis(2*pi/3),
cis (4*pi/3)}⟩
  unfolding image-def by(auto simp: image-def intro:beXI[where x=0]
    beXI[where x=1] beXI[where x=2])
  have ⟨{z. z3 = 1} = {1, cis(2*pi/3), cis (4*pi/3)}⟩
  apply(insert bij-betw-roots-unity[of 3, simplified])
  apply(frule bij-betw-imp-surj-on)
  using f0 by auto
  then show ⟨z3 - 1 = 0 => z = cis (2 * pi / 3) ∨ z = 1 ∨ z = cis (4 * pi
/ 3)⟩
  unfolding bij-betw-def inj-on-def by(auto)
  show ⟨z = cis (2 * pi / 3) ∨ z = 1 ∨ z = cis (4 * pi / 3) => z3 - 1 = 0
  ›
  using cis-2pi cis-multiple-2pi[of 2] by(auto simp:DeMoivre field-simps)
qed

definition root3 where ⟨root3≡cis(2*pi/3)⟩

lemma root3-eq-0:⟨1 + root3 + root32 = 0⟩
proof -
  have ⟨(z::complex)3 - 1 = (z - 1)*(1 + z + z2)⟩ for z
  by(auto simp:field-simps power2-eq-square power3-eq-cube)
  moreover have ⟨root33 - 1 = 0⟩
  using root-unity-3 unfolding root3-def by blast
  moreover have ⟨root3-1 ≠ 0⟩
  by(simp add:cis.code Complex-eq-1 cos-120 root3-def)
  ultimately show ?thesis
  by simp
qed

lemma root-unity-carac:
  assumes h1:⟨a1*a2*a3 = j⟩ ⟨1-a1*a2 ≠ 0 ∧ 1-a2*a3≠0 ∧ 1-a1*a3 ≠ 0⟩
  ⟨j=root3⟩
  and R : ⟨R=(a1 * b2 + b1) / (1 - a1 * a2)⟩ (is ⟨R=?R⟩)
  and P :⟨P=(a2*b3 + b2)/(1-a2*a3)⟩ (is ⟨P=?P⟩)
  and Q:⟨Q=(a3*b1 + b3)/(1-a3*a1)⟩ (is ⟨Q=?Q⟩)
  shows ⟨(a12+a1+1)*b1 + a13*(a22+a2+1)*b2 + a13*a23*(a32+a3+1)*b3

```

```

=
  -j*a1^2*a2*(a1-j)*(a2-j)*(a3-j)*(?R+j*?P+j^2*?Q)
proof -
  have simp-rules-for-eq:⟨1-a1*a2 ≠ 0 ∧ 1-a2*a3 ≠ 0 ∧ 1-a1*a3 ≠ 0⟩ ⟨a1*a2*a3
= j⟩
  ⟨1+j+j^2 = 0⟩ ⟨j*j*j = 1⟩ ⟨(1-a1*a2)*a3 = (a3-j)⟩ ⟨(1-a2*a3)*a1 =
(a1-j)⟩ ⟨(1-a1*a3)*a2 = (a2-j)⟩
  using h1 root3-eq-0 root-unity-3 unfolding root3-def
  by(auto simp:power3-eq-cube mult.left-commute mult.commute right-diff-distrib)

  have y:⟨(-j*a1^2*a2*((1-a2*a3)*a1)*((1-a1*a3)*a2)*((1
-a1*a2)*a3)*
  (a1*b2)+-j*a1^2*a2*((1-a2*a3)*a1)*((1-a1*a3)*a2)*
  ((1-a1*a2)*a3)*b1)/
  (1-a1*a2) = (-j*a1^2*a2*((1-a3*a2)*a1)*((1-a1*a3)*
a2)*
  (a3)*(a1*b2)
  + -j*a1^2*a2*((1-a3*a2)*a1)*((1-a1*a3)*a2)*(a3)*
b1)⟩
  apply(subst add-divide-distrib)
  using simp-rules-for-eq(1)
  by (simp add: mult.commute)
  have y2:⟨(-j*a1^2*a2*((1-a2*a3)*a1)*((1-a1*a3)*a2)*((1
-a1*a2)*a3)*(j*(a2*b3))+
  -j*a1^2*a2*((1-a2*a3)*a1)*((1-a1*a3)*a2)*((1-a1*
a2)*a3)*(j*b2))/
  (1-a2*a3) = (-j*a1^2*a2*(a1)*((1-a1*a3)*a2)*((1-a1
*a2)*a3)*(j*(a2*b3))+
  -j*a1^2*a2*(a1)*((1-a1*a3)*a2)*((1-a1*a2)*a3)*(j*
b2))⟩
  apply(subst add-divide-distrib)
  using simp-rules-for-eq(1) by auto
  have y3:⟨(-j*a1^2*a2*((1-a2*a3)*a1)*((1-a1*a3)*a2)*((1
-a1*a2)*a3)*(j^2*(a3*b1))+
  -j*a1^2*a2*((1-a2*a3)*a1)*((1-a1*a3)*a2)*((1-a1*
a2)*a3)*(j^2*b3))/
  (1-a3*a1) =
  (-j*a1^2*a2*((1-a2*a3)*a1)*(a2)*((1-a1*a2)*a3)*(j^2*
(a3*b1))+
  -j*a1^2*a2*((1-a2*a3)*a1)*(a2)*((1-a1*a2)*a3)*(j^2
*b3))⟩
  apply(subst add-divide-distrib)
  by (smt (verit, best) mult.commute nonzero-mult-div-cancel-left times-divide-eq-left
simp-rules-for-eq(1))
  have ko:⟨a*(b*(c)) = a*b*c⟩ for a b c::complex
  by auto
  have ko':⟨a+(b+(c)) = a+b+c⟩ for a b c::complex
  by auto
  have *:⟨a1*a1*a1*a1*a1*a1*a2*a2*a2*a2*a3*a3*a3*a3*b1
= a1*a1*a2*a3*b1⟩

```

```

    using simp-rules-for-eq by (auto simp add: mult.assoc mult.left-commute)
    have **:⟨a1 * a1 * a1 * a1 * a1 * a2 * a2 * a2 * a2 * a2 * a3 * a3 * a3 *
b3 = a1*a1*a2*a2*b3⟩
    using simp-rules-for-eq by (auto simp add: mult.assoc mult.left-commute)
    have ***:⟨a1 * a1 * a1 * a1 * a1 * a1 * a2 * a2 * a2 * a2 * a2 * a3 * a3 *
a3 * a3 * b3
      = a1*a1*a1*a2*a2*a3*b3⟩
    using simp-rules-for-eq by (auto simp add: mult.assoc mult.left-commute)
    have fin:⟨a1 * b1 + a1 * a1 * a1 * a2 * b2 + a1 * a1 * a1 * a1 * a2 * a2 *
a3 * b2 +
      a1 * a1 * a1 * a2 * a2 * a2 * a3 * b3 + a1 * a1 * a1 * a2 * a2 * a3 * a3
* b1 +
      a1 * a1 * a1 * a1 * a1 * a2 * a2 * a2 * a3 * a3 * b2 + a1 * a1 * a1 * a1 *
a2 * a2 * a2 * a2 *
      a3 * a3 * b3 + a1 * a1 * a2 * a2 * b3 + a1 * a1 * a2 * a3 * b1 = a1*b1
+ j*a1*b1 + j^2*a1*b1 +
      a1*a1*a1*a2*b2 + a1*a1*a1*a2*b2*j + a1*a1*a1*a2*b2*j^2 + a1*a1*a2*a2*b3
+ a1*a1*a2*a2*b3*j +
      a1*a1*a2*a2*b3*j^2⟩
    using simp-rules-for-eq by (auto simp: algebra-simps power2-eq-square power3-eq-cube)
    then have fin':⟨... = (a1*b1 + a1*a1*a1*a2*b2 + a1*a1*a2*a2*b3)*(1+j+j^2)⟩
    by (auto simp: algebra-simps power2-eq-square power3-eq-cube)
    then show graal:⟨(a1^2+a1+1)*b1 + a1^3*(a2^2+a2+1)*b2 + a1^3*a2^3*(a3^2+a3+1)*b3
=
      -j*a1^2*a2*(a1-j)*(a2-j)*(a3-j)*(?R +j*?P +j^2*?Q)⟩
    apply (simp only: simp-rules-for-eq (5-7) [symmetric])
    apply (simp only: y y2 y3 distrib-left distrib-right times-divide-eq-left times-divide-eq-right)

    apply (simp only: simp-rules-for-eq algebra-simps del: mult.assoc mult.commute)

    apply (simp only: ko ko')
    using simp-rules-for-eq apply (auto simp add: power2-eq-square power3-eq-cube
algebra-simps)
    apply (simp only: ko ko') apply (subst *) apply (subst **) apply (subst ***)
  apply (simp)
    apply (simp only: ko ko') apply (subst fin) apply (subst fin')
    using simp-rules-for-eq by fastforce
qed

end
theory Complex-Triangles

```

```

imports Complex-Trigonometry Third-Unity-Root

```

```

begin

```

```

lemma similar-triangles':

```

```

  assumes h: a ≠ 0 ∧ b ≠ 0 ∧ 0 ≠ c ∧ a' ≠ 0 ∧ b' ≠ 0 ∧ c' ≠ 0
  and h1: ∠ (-a) b = ∠ (-a') b' ∧ ∠ (-b') c' = ∠ (-b) c

```

shows $\langle \angle (-c) a = \angle (-c') a' \rangle$
 proof -
 have $\langle \angle (-a) b + \angle (-b) c + \angle (-c) a = \pi \rangle$
 using *angle-sum-triangle'* h by auto
 also have $\langle \angle (-a') b' + \angle (-b') c' + \angle (-c') a' = \pi \rangle$
 using *angle-sum-triangle'* h by auto
 also have $\langle \angle (-a') b' + \angle (-b') c' + \angle (-c) a = \pi \rangle$
 using $\langle \angle (-a) b + \angle (-b) c + \angle (-c) a = \pi \rangle$ h1(1,2) by auto
 then have $\langle \angle (-c) a = \angle (-c') a' \rangle$
 by (smt (verit) $\langle \angle (-a') b' + \angle (-b') c' + \angle (-c') a' = \pi \rangle$ add.inverse-inverse
 ang-vec-bounded
 ang-vec-opposite1 ang-vec-opposite2 ang-vec-opposite-opposite ang-vec-plus-pi1
 ang-vec-plus-pi2
 canon-ang-diff canon-ang-id h neg-0-equal-iff-equal)
 ultimately show ?thesis
 by (metis ang-vec-bounded canon-ang-id)
 qed

lemma *similar-triangles*:
 assumes $h: a \neq b \wedge b \neq c \wedge a \neq c \wedge a' \neq b' \wedge b' \neq c' \wedge c' \neq a'$
 and $h1: \langle \angle (c-a) (b-a) = \angle (c'-a') (b'-a') \rangle \langle \angle (a-b) (c-b) = \angle (a'-b') (c'-b') \rangle$
 shows $\langle \angle (b-c) (a-c) = \angle (b'-c') (a'-c') \rangle$
 proof -
 have $\langle a-b \neq 0 \rangle \langle b-c \neq 0 \rangle \langle a-c \neq 0 \rangle \langle a'-b' \neq 0 \rangle \langle b'-c' \neq 0 \rangle \langle c'-a' \neq 0 \rangle$
 using h by auto
 then show ?thesis
 by (smt (z3) diff-numeral-special(12) h1(1,2) minus-diff-eq similar-triangles')
 qed

lemma *similar-triangles-c*:
 assumes $h: a \neq b \wedge b \neq c \wedge a \neq c \wedge a' \neq b' \wedge b' \neq c' \wedge c' \neq a'$
 and $h1: \langle \text{angle-c } c a b = \text{angle-c } c' a' b' \rangle \langle \text{angle-c } a b c = \text{angle-c } a' b' c' \rangle$
 shows $\langle \text{angle-c } b c a = \text{angle-c } b' c' a' \rangle$
 proof -
 have $\langle a-b \neq 0 \rangle \langle b-c \neq 0 \rangle \langle a-c \neq 0 \rangle \langle a'-b' \neq 0 \rangle \langle b'-c' \neq 0 \rangle \langle c'-a' \neq 0 \rangle$
 using h by auto
 then show ?thesis
 unfolding *angle-c-def*
 by (metis *angle-c-def* h h1(1,2) similar-triangles)
 qed

lemmas *congruent-ctriangleD = congruent-ctriangle.sides congruent-ctriangle.angles*

lemma *congruent-ctriangles-sss*:


```

assumes  $h:a \neq b \wedge b \neq c \wedge a \neq c$ 
and  $h1:\langle cmod (b - a) = cmod (b' - a') \rangle \langle cmod (b - c) = cmod (b' - c') \rangle \langle cmod$ 
 $(c - a) = cmod (c' - a') \rangle$ 
shows  $\langle congruent-ctriangle a b c a' b' c' \rangle$ 
proof -
  {fix  $c b a c' b' a'$ 
    assume  $h:a \neq b \wedge b \neq c \wedge a \neq c$ 
    and  $h1:\langle cmod (b - a) = cmod (b' - a') \rangle \langle cmod (b - c) = cmod (b' - c') \rangle$ 
 $\langle cmod (c - a) = cmod (c' - a') \rangle$ 
    then have  $h':\langle a' \neq b' \wedge b' \neq c' \wedge c' \neq a' \rangle$ 
    by auto
    have  $\langle cos (\angle (c - a) (b - a)) = ((cdist b c)^2 - (cdist a c)^2 - (cdist a b)^2) / (-$ 
 $2 * (cdist a c) * (cdist a b)) \rangle$ 
    using law-of-cosines' h by auto
    moreover have  $\langle cos (\angle (c' - a') (b' - a')) = ((cdist b' c')^2 - (cdist a' c')^2 -$ 
 $(cdist a' b')^2) / (- 2 * (cdist a' c') * (cdist a' b')) \rangle$ 
    using law-of-cosines' h h' by auto
    ultimately have  $f0:\langle cos (\angle (c - a) (b - a)) = cos (\angle (c' - a') (b' - a')) \rangle$ 
    by (metis cdist-def h1(1,2,3) norm-minus-commute)
    have  $\langle \angle (c - a) (b - a) \in \{-pi .. pi\} \rangle \langle \angle (c' - a') (b' - a') \in \{-pi .. pi\} \rangle$ 
    unfolding ang-vec-def
    using canon-ang(1,2) less-eq-real-def by auto
    with  $f0$  have  $\langle \angle (c - a) (b - a) = \angle (c' - a') (b' - a') \vee \angle (c - a) (b - a) = - \angle$ 
 $(c' - a') (b' - a') \rangle$ 
    unfolding ang-vec-def
    by (smt (verit) arccos-cos2 arccos-minus-1 arccos-unique canon-ang(1,2)) } note
dem=this
    then show ?thesis
    by (metis (mono-tags, lifting) angle-c-def cdist-def
congruent-ctriangle-def h h1(1) h1(2) h1(3) norm-minus-commute)
  }
qed

```

lemma *congruent-ctriangleI-sss:*

```

assumes  $h:a \neq b \wedge b \neq c \wedge a \neq c$ 
and  $h1:\langle cdist a b = cdist a' b' \rangle \langle cdist b c = cdist b' c' \rangle \langle cdist a c = cdist a' c' \rangle$ 
shows  $\langle congruent-ctriangle a b c a' b' c' \rangle$ 
using cdist-def congruent-ctriangles-sss h1(1,2,3) h
by (metis dist-commute dist-complex-def)

```

lemmas *congruent-ctriangle-sss = congruent-ctriangleD[OF congruent-ctriangleI-sss]*

lemma *isosceles-triangles:*

```

assumes  $\langle cdist a b = cdist b c \rangle$ 
shows  $\langle angle-c b c a = angle-c b a c \vee angle-c b c a = - angle-c b a c \rangle$ 
by (metis assms cdist-commute cdist-def congruent-ctriangle-sss(14) norm-eq-zero
right-minus-eq)

```

lemma *non-collinear-independant:* $\neg collinear a b c \implies a \neq b \wedge b \neq c \wedge a \neq c$

using *collinear-ex-real* by force

lemma *congruent-ctriangleI-sas*:
 assumes $\langle a1 \neq b1 \wedge b1 \neq c1 \wedge a1 \neq c1 \rangle$
 assumes $h1: \text{cdist } a1 \ b1 = \text{cdist } a2 \ b2$
 assumes $h2: \text{cdist } b1 \ c1 = \text{cdist } b2 \ c2$
 assumes $h3: \text{angle-c } a1 \ b1 \ c1 = \text{angle-c } a2 \ b2 \ c2 \vee \text{angle-c } a1 \ b1 \ c1 = - \text{angle-c } a2 \ b2 \ c2$
 shows *congruent-ctriangle* $a1 \ b1 \ c1 \ a2 \ b2 \ c2$
proof (rule *congruent-ctriangleI-sss*)
 have $h1': \text{cdist } a1 \ b1 = \text{cdist } b2 \ a2$
 using *cdist-commute* $h1$ by auto
 show $\text{cdist } a1 \ c1 = \text{cdist } a2 \ c2$
proof (rule *power2-eq-imp-eq*)
 show $(\text{cdist } a1 \ c1)^2 = (\text{cdist } a2 \ c2)^2$
 apply (insert *law-of-cosines*[of $a1 \ c1 \ b1$] *law-of-cosines*[of $a2 \ c2 \ b2$])
 apply (subst (asm) (1 2) $h2$) apply (subst (asm) (1 2) $h1'$ [symmetric])
 using *assms* **unfolding** *angle-c-def*
 by (smt (verit) *congruent-ctriangle-sss*(7) *angle-c-commute* *angle-c-def* *cos-minus*)
qed *simp-all*
 show $a1 \neq b1 \wedge b1 \neq c1 \wedge a1 \neq c1$
 $\text{cdist } a1 \ b1 = \text{cdist } a2 \ b2$
 $\text{cdist } b1 \ c1 = \text{cdist } b2 \ c2$ using *assms* *cdist-commute* by auto
qed

lemmas *congruent-ctriangle-sas* = *congruent-ctriangleD*[OF *congruent-ctriangleI-sas*]

lemma *congruent-ctriangleI-aas*:
 assumes $h1: \text{angle-c } a1 \ b1 \ c1 = \text{angle-c } a2 \ b2 \ c2$
 assumes $h2: \text{angle-c } b1 \ c1 \ a1 = \text{angle-c } b2 \ c2 \ a2$
 assumes $h3: \text{cdist } a1 \ b1 = \text{cdist } a2 \ b2$
 assumes $h4: \neg \text{collinear } a1 \ b1 \ c1 \ \neg \text{collinear } a2 \ b2 \ c2$
 shows *congruent-ctriangle* $a1 \ b1 \ c1 \ a2 \ b2 \ c2$
proof (rule *congruent-ctriangleI-sas*)
 from $h4$ have $\text{neg}: a1 \neq b1$ **unfolding** *collinear-def* by auto
 with *assms*(3) have $\text{neg}': a2 \neq b2$ by auto
 have $h0: \langle a1 \neq b1 \rangle \langle b1 \neq c1 \rangle \langle a1 \neq c1 \rangle$
 using *assms*(4) *non-collinear-independant* by auto
 have $h0': \langle a2 \neq b2 \rangle \langle b2 \neq c2 \rangle \langle a2 \neq c2 \rangle$
 using *assms*(5) *non-collinear-independant* by auto
 have $A: \text{angle-c } c1 \ a1 \ b1 = \text{angle-c } c2 \ a2 \ b2$ using neg neg' *assms*
 apply (insert *angle-sum-triangle-c*[of $a1 \ b1 \ c1$] *angle-sum-triangle-c*[of $a2 \ b2 \ c2$])
 by (metis $h0$ *assms*(5) $h1 \ h2$ *non-collinear-independant* *similar-triangles-c*)
 have $\langle -\angle (b1 - c1) (a1 - c1) = -\angle (b2 - c2) (a2 - c2) \rangle$

```

    using h2 unfolding angle-c-def by auto
  then have A1:⟨angle-c a1 c1 b1 = angle-c a2 c2 b2⟩
    using h2 unfolding angle-c-def
    apply(cases ⟨angle-c a1 c1 b1 = pi⟩)
    apply (metis ang-vec-def angle-c-def canon-ang-uminus-pi minus-diff-eq)
    by (metis ang-vec-sym ang-vec-sym-pi)
  have ⟨∠ (b1 - a1) (c1 - a1) = ∠ (b2 - a2) (c2 - a2)⟩
    by (metis ang-vec-sym ang-vec-sym-pi angle-c-def A)
  have sine1:⟨ sin (angle-c a1 c1 b1) * cdist c1 b1 = sin (angle-c b1 a1 c1) * cdist
a1 b1⟩
    by (metis h0 law-of-sines)
  have sine2:⟨sin (angle-c a2 c2 b2) * cdist c2 b2 = sin (angle-c b2 a2 c2) * cdist
a2 b2⟩
    by (metis assms(5) law-of-sines non-collinear-independant)
  have ⟨¬collinear a2 b2 c2 ⟹ sin (angle-c a2 c2 b2) ≠ 0⟩
    unfolding collinear-def angle-c-def
    using angle-c-def collinear-sin-neq-0 h4(2) by presburger
  have h3':⟨cdist b1 a1 = cdist b2 a2⟩
    using h3 cdist-commute by auto
  have A2:⟨angle-c b1 a1 c1 = angle-c b2 a2 c2⟩
    using ⟨∠ (b1 - a1) (c1 - a1) = ∠ (b2 - a2) (c2 - a2)⟩ angle-c-def by auto
  have ⟨ cdist c1 b1 = sin (angle-c b2 a2 c2) * cdist a2 b2 / sin (angle-c a2 c2
b2) ∧
    cdist c2 b2 = sin (angle-c b2 a2 c2) * cdist a2 b2 / sin (angle-c a2 c2 b2) ⟩
    apply(cases ⟨ sin (angle-c a2 c2 b2) = 0 ⟩)
    apply (simp add:
      ⟨¬ Elementary-Complex-Geometry.collinear a2 b2 c2 ⟹ sin (angle-c a2 c2
b2) ≠ 0⟩
      ⟨¬ Elementary-Complex-Geometry.collinear a2 b2 c2⟩)
  by (metis A1 A2 h3 nonzero-mult-div-cancel-left sine1 sine2)
  then show cdist b1 c1 = cdist b2 c2
    using cdist-commute by auto
  show a1 ≠ b1 ∧ b1 ≠ c1 ∧ a1 ≠ c1
    cdist a1 b1 = cdist a2 b2
    angle-c a1 b1 c1 = angle-c a2 b2 c2 ∨ angle-c a1 b1 c1 = - angle-c a2 b2 c2
    using assms h0 neq by auto
qed

```

lemmas congruent-ctriangle-aas = congruent-ctriangleD[OF congruent-ctriangleI-aas]

lemma congruent-ctriangleI-asa:

```

  assumes angle-c a1 b1 c1 = angle-c a2 b2 c2
  assumes cdist a1 b1 = cdist a2 b2
  assumes h0:angle-c b1 a1 c1 = angle-c b2 a2 c2
  assumes h4:¬collinear a1 b1 c1 ¬collinear a2 b2 c2
  shows congruent-ctriangle a1 b1 c1 a2 b2 c2
  proof (rule congruent-ctriangleI-aas)
    from assms have neg: a1 ≠ b1 a2 ≠ b2 unfolding collinear-def by auto
  qed

```

```

show angle-c b1 c1 a1 = angle-c b2 c2 a2
  apply (rule similar-triangles-c)
  using assms(4,5) non-collinear-independant apply auto[1]
  using h0 unfolding angle-c-def
  apply (metis ang-vec-sym ang-vec-sym-pi)
  using angle-c-def assms(1) by auto
qed fact+

```

```

lemmas congruent-ctriangle-asa = congruent-ctriangleD[OF congruent-ctriangleI-asa]

```

```

lemma orientation-respect-letter-order:
  assumes angle-c a b c = angle-c b a c  $\neg$  collinear a b c
  shows False
  by (smt (verit, ccfv-threshold) angle-c-commute assms(1)
      assms(2) collinear-sin-neq-0 collinear-sym1 similar-triangles-c sin-pi sin-zero)

```

```

lemma isoscele-iff-pr-cis-qr:
  assumes h': $\langle q \neq r \rangle$ 
  shows  $\langle \text{cdist } q \ r = \text{cdist } p \ r \rangle \longleftrightarrow (p-r) = \text{cis}(\text{angle-c } q \ r \ p) * (q-r) \rangle$ 
proof (rule iffI)
  assume h: $\langle \text{cdist } q \ r = \text{cdist } p \ r \rangle$ 
  then have f0: $\langle \text{cmod}(p-r) = \text{cmod}(q-r) \rangle$ 
    by (simp add: norm-minus-commute)
  have  $\langle (p-r)/(q-r) = (p-r)*\text{cnj}(q-r)/\text{cmod}(q-r)^2 \rangle$ 
    using complex-div-cnj by blast
  have f1: $\langle \text{Re}((p-r)*\text{cnj}(q-r)) = \text{cmod}(p-r)*\text{cmod}(q-r)*\cos(\text{angle-c } q \ r \ p) \rangle$ 
    by (metis (no-types, lifting) ang-cos cdist-commute cdist-def h(1))
  have f2: $\langle \text{Im}((p-r)*\text{cnj}(q-r)) = \text{cmod}(p-r)*\text{cmod}(q-r)*\sin(\text{angle-c } q \ r \ p) \rangle$ 
    using ang-sin h(1) by auto
  then have  $\langle (p-r)*\text{cnj}(q-r)/\text{cmod}(q-r)^2 = \text{cmod}(p-r)*\text{cmod}(q-r)*\text{cis}(\text{angle-c } q \ r \ p)/\text{cmod}(q-r)^2 \rangle$ 
    apply (intro complex-eqI)
    using f1 f2 by auto
  then have f3: $\langle (p-r)/(q-r) = \text{cmod}(p-r)*\text{cmod}(q-r)*\text{cis}(\text{angle-c } q \ r \ p)/\text{cmod}(q-r)^2 \rangle$ 
    by (simp add:  $\langle (p-r) / (q-r) = (p-r) * \text{cnj}(q-r) / \text{cor}((\text{cmod}(q-r))^2) \rangle$ )
  have  $\langle (p-r)/(q-r) = \text{cis}(\text{angle-c } q \ r \ p) \rangle$ 
    apply (subst f3) apply (subst f0) by (simp add: h h' power2-eq-square)
  then show  $\langle p-r = \text{cis}(\text{angle-c } q \ r \ p) * (q-r) \rangle$ 
    using divide-eq-eq h h' by force
next
  assume  $\langle p-r = \text{cis}(\text{angle-c } q \ r \ p) * (q-r) \rangle$ 
  have  $\langle (p-r)/(q-r) = (p-r)*\text{cnj}(q-r)/\text{cmod}(q-r)^2 \rangle$ 
    using complex-div-cnj by blast
  have f1: $\langle \text{Re}((p-r)*\text{cnj}(q-r)) = \text{cmod}(p-r)*\text{cmod}(q-r)*\cos(\text{angle-c } q \ r \ p) \rangle$ 
    using ang-cos by fastforce

```

```

have f2:  $\langle \text{Im } ((p-r) * \text{cnj}(q-r)) = \text{cmod}(p-r) * \text{cmod}(q-r) * \sin(\text{angle-c } q \ r \ p) \rangle$ 
using ang-sin by fastforce
then have f4:  $\langle (p-r) * \text{cnj}(q-r) / \text{cmod}(q-r)^2 = \text{cmod}(p-r) * \text{cmod}(q-r) * \text{cis}(\text{angle-c } q \ r \ p) / \text{cmod}(q-r)^2 \rangle$ 
apply(intro complex-eqI)
using f1 f2 by auto
then have f3:  $\langle (p-r) / (q-r) = \text{cmod}(p-r) * \text{cmod}(q-r) * \text{cis}(\text{angle-c } q \ r \ p) / \text{cmod}(q-r)^2 \rangle$ 
by (simp add:  $\langle (p-r) / (q-r) = (p-r) * \text{cnj}(q-r) / \text{cor}((\text{cmod}(q-r))^2) \rangle$ )
have  $\langle \text{cmod}(p-r) * \text{cmod}(q-r) * \text{cis}(\text{angle-c } q \ r \ p) / \text{cmod}(q-r)^2 = \text{cis}(\text{angle-c } q \ r \ p) \rangle$ 
using  $\langle p-r = \text{cis}(\text{angle-c } q \ r \ p) * (q-r) \rangle \langle q \neq r \rangle$  f3 by auto
then have  $\langle \text{cmod}(p-r) * \text{cmod}(q-r) / \text{cmod}(q-r)^2 = 1 \rangle$ 
by (metis f4  $\langle q \neq r \rangle$  cmod-power2 complex-mod-cn timer-complex-mod-mult-cn timer-complex-mult-cn timer-eq-iff-diff-eq-0 nonzero-norm-divide norm-cis norm-eq-zero norm-mult power-not-zero)
then show  $\langle \text{cdist } q \ r = \text{cdist } p \ r \rangle$ 
by (metis cdist-commute cdist-def eq-divide-eq-1 mult.commute power2-eq-square real-divide-square-eq)
qed

```

lemma *equilateral-imp-pi3*:

```

assumes  $\langle q \neq r \rangle \text{cdist } q \ r = \text{cdist } p \ r \text{cdist } p \ r = \text{cdist } p \ q$ 
shows  $\lfloor (\text{angle-c } q \ r \ p) \rfloor = \pi/3 \vee \lfloor (\text{angle-c } q \ r \ p) \rfloor = -\pi/3$ 
 $(\text{angle-c } q \ r \ p) = (\text{angle-c } p \ q \ r) \wedge (\text{angle-c } q \ r \ p) = (\text{angle-c } r \ p \ q)$ 
proof -
have f0:  $\langle q \neq r \wedge r \neq p \wedge p \neq q \rangle$  using assms by(auto)
have  $\langle \text{angle-c } q \ r \ p \neq 0 \wedge \text{angle-c } q \ r \ p \neq \pi \rangle$ 
by (smt (verit, best) Re-complex-of-real angle-c-commute angle-c-commute-pi
assms(1) assms(2)
assms(3) cdist-commute cdist-def congruent-ctriangle-sss(20) cor-cmod-real
isoscele-iff-pr-cis-qr
minus-complex.simps(1) minus-complex.simps(2) minus-diff-eq mult-minus-right
norm-eq-zero right-minus-eq)
then have f1:  $\langle \neg \text{collinear } q \ r \ p \rangle$ 
using collinear-angle f0 by blast
then have f1':  $\langle \neg \text{collinear } p \ q \ r \rangle$ 
by (simp add: collinear-sym1 collinear-sym2)
have f12:  $\langle (\text{angle-c } q \ r \ p) = \text{angle-c } r \ p \ q \vee (\text{angle-c } q \ r \ p) = -\text{angle-c } r \ p \ q \rangle$ 
apply(rule congruent-ctriangle-sss)
using f0 assms by(auto simp:norm-minus-commute)
have f4:  $\langle (\text{angle-c } q \ r \ p) = \text{angle-c } p \ q \ r \vee (\text{angle-c } q \ r \ p) = -\text{angle-c } p \ q \ r \rangle$ 
apply(rule congruent-ctriangle-sss)
using f0 assms by(auto simp:norm-minus-commute)
have f5:  $\langle (\text{angle-c } q \ r \ p) \neq \pi \rangle$ 
using f1
by (metis angle-c-commute canon-ang-sin canon-ang-uminus-pi collinear-sin-neq-0
f1' pi-canonical)

```

```

    sin-pi)
  from f4 have ⟨ |(angle-c q r p) + (angle-c r p q) + (angle-c p q r)| = pi ⟩
    using angle-sum-triangle-c[of p q r] f0 by auto
  have ⟨ ∀ x ∈ {-pi..0}. 3*x ≠ pi ⟩
    by (metis atLeastAtMost-iff dual-order.trans mult-eq-0-iff mult-less-cancel-right2
numeral-le-one-iff
    pi-ge-zero pi-neq-zero semiring-norm(70) verit-comp-simplify1(3) verit-la-disequality)
  also have the-x:⟨ ∃ !x ∈ {-pi<..pi}. 3*x = pi ⟩
    by(auto intro:exI[where x=⟨pi/3⟩])
  moreover have ⟨ (angle-c x y z) ∈ {-pi<..pi} ⟩ for x y z
    using ang-c-in by auto
  have ⟨ x ∈ {-3*pi<..3*pi} ∧ x-2*pi = pi ⟹ ( x = 3*pi) ⟩ for x k::real
    by(auto simp:)
  have ⟨ x ∈ {-3*pi<..3*pi} ∧ x-2*k*pi = pi ∧ ( k ≥ 2 ∨ k ≤ - 2) ⟹ False ⟩ for
x and k::int
  proof -
    assume h:⟨ x ∈ {-3*pi<..3*pi} ∧ x-2*k*pi = pi ∧ ( k ≥ 2 ∨ k ≤ -2) ⟩
    then have f0:⟨ r = 2 ⟹ x - 2* r*pi ≤ -pi ⟩ for r by auto
    also have f1:⟨ k > 2 ⟹ x - 2* k*pi < x-2*2*pi ⟩
      by(auto)
    have f3:⟨ x-2*2*pi ≤ -pi ⟩
      using f0[of 2] by(auto simp:abs-real-def)
    then have f2:⟨ k > 2 ⟹ x - 2* k*pi < -pi ⟩
      using f1 by argo
    ultimately have f3:⟨ k ≥ 2 ⟹ x-2* k*pi ≤ -pi ⟩
      apply(cases ⟨abs k=2⟩) using f0 by(auto)
    have f0':⟨ r=-2 ⟹ x - 2*r*pi > pi ⟩ for r ::int
      using h by(auto)
    have ⟨ x+2*2*pi > pi ⟩ using h by auto
    also have ⟨ k < -2 ⟹ x - 2*k*pi > x - 2*(-2)*pi ⟩
      using f0'[of 2]
    proof -
      assume a1: k < - 2
      have f2: pi < x - real-of-int (- 4) * pi
        using f0' by force
      have k + k ≤ - 4
        using a1 by force
      then show ?thesis
        using f2 by (metis (no-types) diff-left-mono h mult.commute mult-2-right
of-int-le-iff ordered-comm-semiring-class.comm-mult-left-mono pi-ge-zero verit-comp-simplify1(3))
    qed
    ultimately have f4:⟨ k ≤ -2 ⟹ x-2* k*pi > pi ⟩
      apply(cases ⟨k=-2⟩) using h
      ⟨ k < - 2 ⟹ x - 2 * - 2 * pi < x - real-of-int (2 * k) * pi ⟩ by auto
    show False
      using f3 f4 h by force
  qed
  then have possible-k:⟨ x ∈ {-3*pi<..3*pi} ∧ x-2*k*pi = pi ⟹ k=-1 ∨ k=1
  ∨ k=0 ⟩ for k::int and x

```

```

    by force
    have ang-3-in:⟨y ∈ {−3*pi<..3*pi} ∧ |y| = pi ⟷ y=pi ∨ y = 3*pi ∨ y=−pi⟩
  for y
  proof(rule iffI)
    assume hyp:⟨y ∈ {−3 * pi<..3 * pi} ∧ |y| = pi⟩
    have ⟨∃ k. y − 2*k*pi = pi⟩
      by (metis add-diff-cancel-left' diff-add-cancel divide-self-if field-sum-of-halves
        mult.commute mult-1 mult-2
        pi-neq-zero times-divide-eq-right)
    with hyp show ⟨y = pi ∨ y=3*pi ∨ y= −pi⟩
      using hyp possible-k canon-ang-pi-3pi[of y] canon-ang-id canon-ang-uminus-pi

    by (smt (verit, del-insts) canon-ang-minus-3pi-minus-pi greaterThanAtMost-iff)
  next
    assume ⟨y = pi ∨ y = 3 * pi ∨ y = − pi⟩
    then show ⟨y ∈ {−3 * pi<..3 * pi} ∧ |y| = pi⟩
      by(auto simp:canon-ang-uminus-pi canon-ang-pi-3pi)
  qed
  ultimately have ⟨|3*(angle-c q r p)| = pi ⟹ angle-c q r p = −pi/3 ∨ angle-c
q r p = pi/3⟩
  proof −
    assume ⟨|3*(angle-c q r p)| = pi⟩
    have ⟨3*(angle-c q r p) ∈ {−3*pi<..3*pi}⟩
      using ang-c-in by auto
    then have f0:⟨3*(angle-c q r p) = −pi ∨ 3*(angle-c q r p) = pi ∨ 3*(angle-c
q r p) = 3*pi⟩
      using ang-3-in ⟨|3*(angle-c q r p)| = pi⟩ by blast
    then have ⟨3*(angle-c q r p) = 3*pi ⟹ False⟩
      using collinear-sin-neq-pi by (simp add: f5)
    then show ⟨angle-c q r p = −pi/3 ∨ angle-c q r p = pi/3⟩
      using f0 by auto
  qed
  have f2:⟨|angle-c x y z| = angle-c x y z⟩ for x y z::complex
    using ang-vec-bounded angle-c-def canon-ang-id by fastforce
  then have ⟨|3*(angle-c q r p)| = pi⟩
    apply (cases ⟨(angle-c q r p) = − angle-c r p q⟩)
    apply(simp add:f1 f2)
    apply (metis add.commute add.right-inverse add.right-neutral angle-sum-triangle-c
      canon-ang-uminus-pi f0 f2 f4 f5)
    by (smt (verit) ⟨|angle-c q r p + angle-c r p q + angle-c p q r| = pi⟩
      ⟨angle-c q r p = angle-c p q r ∨ angle-c q r p = − angle-c p q r⟩
      ⟨angle-c q r p = angle-c r p q ∨ angle-c q r p = − angle-c r p q⟩
      canon-ang-uminus-pi
      f2)
  then have ⟨angle-c q r p = −pi/3 ∨ angle-c q r p = pi/3⟩
    using ⟨|3 * angle-c q r p| = pi ⟹ angle-c q r p = − pi / 3 ∨ angle-c q r p
= pi / 3⟩ by auto
  then show ⟨|angle-c q r p| = pi / 3 ∨ |angle-c q r p| = − pi / 3⟩
    using f2 by auto

```

then show $\langle \text{angle-c } q \ r \ p = \text{angle-c } p \ q \ r \wedge \text{angle-c } q \ r \ p = \text{angle-c } r \ p \ q \rangle$
by (smt (verit, del-insts) $\langle \text{angle-c } q \ r \ p + \text{angle-c } r \ p \ q + \text{angle-c } p \ q \ r \rangle = \pi$)

 $f12 \text{ collinear-sin-neq-0 collinear-sym1 } f1' \ f2 \ f4 \ f5 \ \sin\pi$
qed

lemma isosceles-triangle-converse:
assumes $\text{angle-c } a \ b \ c = \text{angle-c } c \ a \ b \neg \text{collinear } a \ b \ c$
shows $\text{dist } a \ c = \text{dist } b \ c$
by (metis assms(1) assms(2) cdist-def collinear-sin-neq-0 collinear-sym1 congruent-ctriangle-asa(8)
congruent-ctriangle-asa(9) dist-complex-def law-of-sines mult-cancel-left non-collinear-independant)

lemma pi3-imp-equilateral:
assumes $\langle q \neq r \rangle \langle p \neq q \rangle \langle r \neq p \rangle$
and $\langle (\text{angle-c } q \ p \ r) = \pi/3 \vee (\text{angle-c } q \ p \ r) = -\pi/3 \rangle$
and $\langle (\text{angle-c } q \ p \ r) = (\text{angle-c } r \ q \ p) \rangle$
and $\langle (\text{angle-c } q \ p \ r) = (\text{angle-c } p \ r \ q) \rangle$
shows $\langle \text{cdist } p \ r = \text{cdist } q \ r \wedge \text{cdist } p \ r = \text{cdist } p \ q \rangle$
proof(safe)
have $f0: \langle \neg \text{collinear } q \ p \ r \rangle$
using assms(1-3) **unfolding** collinear-def
by (smt (verit, ccfv-SIG) arccos-one-half arcsin-one-half arcsin-plus-arccos
assms(4)
collinear-angle collinear-def divide-le-eq-1-pos divide-pos-pos minus-divide-left
pi-def pi-gt-zero pi-half)
then show $\langle \text{cdist } p \ r = \text{cdist } q \ r \rangle$
using isosceles-triangle-converse[of $q \ p \ r$] assms
by (simp add: dist-complex-def norm-minus-commute)
show $\langle \text{cdist } p \ r = \text{cdist } p \ q \rangle$
using $f0$ isosceles-triangle-converse[of $p \ r \ q$] assms
by (metis $\langle \text{cdist } p \ r = \text{cdist } q \ r \rangle$ cdist-def collinear-iff dist-commute dist-norm)
qed

lemma pi3-isoscele-imp-equilateral:
assumes $\langle q \neq r \rangle \langle p \neq q \rangle \text{cdist } q \ r = \text{cdist } p \ r$
and $\lfloor (\text{angle-c } q \ p \ r) \rfloor = \pi/3 \vee \lfloor (\text{angle-c } q \ p \ r) \rfloor = -\pi/3$
shows $\langle \text{cdist } p \ q = \text{cdist } r \ q \rangle$
proof(-)
have $f: \langle \text{angle-c } p \ q \ r = (\text{angle-c } r \ p \ q) \rangle$
by (smt (verit, best) angle-c-commute angle-c-commute-pi assms(2) assms(3)
cdist-commute
collinear-angle isosceles-triangles orientation-respect-letter-order)
have $f0: \langle r \neq p \rangle$
using assms(1) assms(3) **by** force
then have $\langle \text{cdist } q \ p = \text{cdist } r \ q \rangle$
by (smt (verit) congruent-ctriangle-sss(19) add-divide-distrib ang-vec-bounded)

angle-c-commute angle-c-def $\text{angle-sum-triangle-c}$ $\text{arccos-minus-one-half ar-ccos-one-half arcsin-minus-one-half}$
 $\text{arcsin-one-half arcsin-plus-arccos}$ $\text{assms}(2)$ $\text{assms}(3)$ $\text{assms}(4)$ canon-ang-id
 $\text{canon-ang-minus-3pi-minus-pi}$
 cdist-commute $\text{field-sum-of-halves}$ law-of-sines mult-cancel-right $\text{pi3-imp-equilateral}$
 sin-gt-zero sin-periodic-pi
then show $\langle \text{cdist } p \text{ } q = \text{cdist } r \text{ } q \rangle$
using cdist-commute **by force**
qed

lemma $\text{pi3-isoscele-imp-equilateral'}$:
assumes $\langle q \neq r \rangle$ $\langle p \neq q \rangle$ $\text{cdist } q \text{ } r = \text{cdist } p \text{ } r$
and $\lfloor (\text{angle-c } q \text{ } p \text{ } r) \rfloor = \text{pi}/3 \vee \lfloor (\text{angle-c } q \text{ } p \text{ } r) \rfloor = -\text{pi}/3$
shows $\langle \text{cdist } p \text{ } r = \text{cdist } p \text{ } q \rangle$
by ($\text{metis assms}(2)$ $\text{assms}(3)$ $\text{assms}(4)$ cdist-commute $\text{pi3-isoscele-imp-equilateral}$)

lemma $\text{equilateral-characterization}$: $\langle q \neq r \implies (\text{cdist } q \text{ } r = \text{cdist } p \text{ } r \wedge \text{cdist } p \text{ } r = \text{cdist } p \text{ } q) \iff ((p-r) = \text{cis}(\text{pi}/3) * (q-r) \vee (p-r) = \text{cis}(-\text{pi}/3) * (q-r)) \rangle$
proof(rule iffI)
assume h : $\langle q \neq r \rangle$ $\langle \text{cdist } q \text{ } r = \text{cdist } p \text{ } r \wedge \text{cdist } p \text{ } r = \text{cdist } p \text{ } q \rangle$
then have $f1$: $\langle p - r = \text{cis}(\text{angle-c } q \text{ } r \text{ } p) * (q - r) \vee p - r = \text{cis}(-\text{angle-c } q \text{ } r \text{ } p) * (q - r) \rangle$
using $\text{isoscele-iff-pr-cis-qr}$ **by blast**
have $f0$: $\langle q \neq r \wedge r \neq p \wedge p \neq q \rangle$ **using** h **by auto**
have $\langle \text{angle-c } q \text{ } r \text{ } p = \text{pi}/3 \vee \text{angle-c } q \text{ } r \text{ } p = -\text{pi}/3 \rangle$
using $\text{equilateral-imp-pi3}(1)$ **[of** $q \text{ } r \text{ } p$ **]**
by ($\text{metis ang-vec-bounded}$ angle-c-def canon-ang-id $h(1)$ $h(2)$)
then show $\langle p - r = \text{cis}(\text{pi} / 3) * (q - r) \vee p - r = \text{cis}(-\text{pi} / 3) * (q - r) \rangle$
using $f1$
by ($\text{metis add.inverse-inverse minus-divide-left}$)
next

assume h : $\langle q \neq r \rangle$ $\langle p - r = \text{cis}(\text{pi} / 3) * (q - r) \vee p - r = \text{cis}(-\text{pi} / 3) * (q - r) \rangle$
have $\langle (p-r)/(q-r) = (p-r)*\text{cnj}(q-r)/\text{cmod}(q-r)^2 \rangle$
using complex-div-cnj **by blast**
have $f1$: $\langle \text{Re}((p-r)*\text{cnj}(q-r)) = \text{cmod}(p-r)*\text{cmod}(q-r)*\cos(\text{angle-c } q \text{ } r \text{ } p) \rangle$
using ang-cos **by force**
have $f2$: $\langle \text{Im}((p-r)*\text{cnj}(q-r)) = \text{cmod}(p-r)*\text{cmod}(q-r)*\sin(\text{angle-c } q \text{ } r \text{ } p) \rangle$
using ang-sin $h(1)$ **by auto**
then have $\langle (p-r)*\text{cnj}(q-r)/\text{cmod}(q-r)^2 = \text{cmod}(p-r)*\text{cmod}(q-r)*\text{cis}(\text{angle-c } q \text{ } r \text{ } p)/\text{cmod}(q-r)^2 \rangle$
apply(intro complex-eqI)
using $f1$ $f2$ **by auto**
then have $f3$: $\langle (p-r)/(q-r) = \text{cmod}(p-r)*\text{cmod}(q-r)*\text{cis}(\text{angle-c } q \text{ } r \text{ } p)/\text{cmod}(q-r)^2 \rangle$
by (simp add: $\langle (p - r) / (q - r) = (p - r) * \text{cnj}(q - r) / \text{cor}((\text{cmod}(q -$

```

r))2)
  have ⟨(p-r)/(q-r) = cis (angle-c q r p)⟩
  apply(subst f3)
  by (smt (z3) ⟨(p - r) * cnj (q - r) / cor ((cmod (q - r))2) = cor (cmod (p
- r) * cmod (q - r))
* cis (angle-c q r p) / cor ((cmod (q - r))2)⟩ cis-divide cis-mult cis-neq-zero
cis-zero
  cmod-cor-divide complex-mod-cn timeriffdiff-eq-0 f3 h(1) h(2) nonzero-mult-divide-mult-cancel-left

  nonzero-mult-divide-mult-cancel-right norm-cis norm-mult numeral-One
of-real-1 of-real-divide
  power2-less-0 times-divide-eq-left)
  then have fpi:⟨angle-c q r p = pi/3 ∨ angle-c q r p = -pi/3⟩
  using h ang-c-in
  by (smt (verit, best) add-divide-distrib ang-vec-bounded angle-c-def arccos-one-half

  arcsin-one-half arcsin-plus-arccos cis-inj diff-eq-diff-eq diff-numeral-special(9)

  divide-nonneg-pos field-sum-of-halves nonzero-divide-eq-eq)
  then have ff:⟨cdist p r = cdist q r⟩
  by (metis ⟨(p - r) / (q - r) = cis (angle-c q r p)⟩ h(1) isoscele-iff-pr-cis-qr
nonzero-divide-eq-eq right-minus-eq)
  then have ⟨angle-c p q r = angle-c q p r ∨ angle-c p q r = -angle-c q p r ⟩
  by (metis congruent-ctriangle-sss(13) cdist-commute cdist-def norm-eq-zero
right-minus-eq)
  then have f10:⟨|angle-c p q r + angle-c q r p + angle-c r p q| = pi⟩
  by (metis ⟨angle-c q r p = pi / 3 ∨ angle-c q r p = - pi / 3⟩ angle-c-neq0
angle-sum-triangle-c
  divide-eq-0-iff h(1) neg-equal-0-iff-equal pi-neq-zero zero-neq-numeral)
  then have ⟨|angle-c q p r| = pi / 3 ∨ |angle-c q p r| = - pi / 3⟩
  by (smt (verit) ⟨angle-c p q r = angle-c q p r ∨ angle-c p q r = - angle-c q p
r⟩
  add-divide-distrib ang-pos-pos ang-vec-bounded ang-vec-sym angle-c-def arc-
cos-minus-one-half
  arccos-one-half arcsin-minus-one-half arcsin-one-half arcsin-plus-arccos
canon-ang-id
  canon-ang-minus-3pi-minus-pi canon-ang-pi-3pi field-sum-of-halves fpi)
  show ⟨cdist q r = cdist p r ∧ cdist p r = cdist p q⟩
  apply(rule conjI)
  using ff apply(auto simp: norm-minus-commute)[1]
  apply(rule pi3-isoscele-imp-equilateral)
  using h(1) apply auto[1]
  apply (metis ⟨angle-c q r p = pi / 3 ∨ angle-c q r p = - pi / 3⟩
angle-c-neq0 divide-eq-0-iff neg-equal-0-iff-equal pi-neq-zero zero-neq-numeral)
  using ff apply presburger
  using ⟨|angle-c q p r| = pi / 3 ∨ |angle-c q p r| = - pi / 3⟩ by blast
qed

lemmas equilateral-imp-prcispi3 = equilateral-characterization[THEN iffD1]

```

lemmas *prcispi3-imp-equilateral* = *equilateral-characterization*[*THEN iffD2*]

lemma *equilateral-characterization-neg*:

```

  fixes p q r::complex
  assumes h1:⟨p≠r⟩
  shows ⟨(cdist p r = cdist p q ∧ cdist p q = cdist q r ∧ angle-c q r p = -pi/3)
  ⟷ p + root3 * q + root3 ^ 2 * r = 0⟩
proof(rule iffI)
  assume h:⟨cdist p r = cdist p q ∧ cdist p q = cdist q r ∧ angle-c q r p = -pi/3⟩
  then have ⟨p-r = cis(-pi/3)*(q-r)⟩
    using h1 isoscele-iff-pr-cis-qr[of q r p]
    by force
  with h have ⟨p-r+cis(2*pi/3)*(q-r) = 0⟩
    apply(simp add:cis.code)
    apply(subst sin-120)
    apply(intro complex-eqI)
    using cos-120 cos-60 sin-120 sin-60
    by(auto simp add:field-simps)
  then have ⟨p + root3*q - (root3+1)*r = 0⟩ by(auto simp:field-simps root3-def)
  then show ⟨p + root3 * q + root3^2 * r = 0⟩
    by (metis (no-types, lifting) add.commute eq-iff-diff-eq-0 left-diff-distrib
    ring-class.ring-distribs(2) root3-eq-0)
next
  assume h:⟨p + root3 * q + root3^2 * r = 0⟩
  then have ⟨p + root3*q - (root3+1)*r = 0⟩
    by (metis add.commute add-diff-cancel-left diff-numeral-special(9) left-diff-distrib
    ring-class.ring-distribs(2)
    root3-eq-0)
  then have ⟨p-r+root3*(q-r) = 0⟩
    by(auto simp:field-simps)
  have ⟨-root3 = cis(-pi/3)⟩
    using cos-120 cos-60 sin-120 sin-60
    by (auto intro:complex-eqI simp: root3-def)
  then have f1':⟨p - r = cis(-pi/3)*(q-r)⟩
    by (metis ⟨p - r + root3 * (q - r) = 0⟩ eq-neg-iff-add-eq-0 mult-minus-left)
  then have f1:⟨p - r = cis(pi/3)*(q-r) ∨ p - r = cis(-pi/3)*(q-r)⟩
    by auto
  then have f2:⟨cmod (p-r) = cmod (q-r)⟩
    by(auto simp:norm-mult)
  then have f3:⟨dist p r = dist q r⟩
    by (simp add: dist-complex-def)
  have ang-dem:⟨angle-c q r p = -pi/3⟩
proof -
  have *:⟨q≠r⟩
    using f2 h1 by auto
  then have **:⟨(p - r) / (q-r) = cis(-pi/3)⟩
    using f1' by auto
  have ***:⟨angle-c q r p = Arg ((p - r) / (q-r))⟩

```

```

    by (metis ang-vec-def angle-c-def arg-div f1' h1 mult-zero-right right-minus-eq)
  then have  $\langle \text{Arg } ((p - r) / (q - r)) = \lfloor -\pi/3 \rfloor \rangle$ 
    by (simp add: arg-cis **)
  then show ?thesis using * ** ***
    using canon-ang-id pi-ge-two by force
qed
have  $\langle \text{cmod } (q - p) = \text{cmod } (r - q) \rangle$ 
  using f2
  by (metis cdist-def dist-eq-0-iff equilateral-characterization f1' f3)
then show  $\langle \text{cdist } p \ r = \text{cdist } p \ q \wedge \text{cdist } p \ q = \text{cdist } q \ r \wedge \text{angle-c } q \ r \ p = -\pi/3 \rangle$ 

  using equilateral-characterization[THEN iffD2, of q r p, OF - f1]
  using f2 ang-dem by(auto simp: f2 field-simps intro!)
qed

```

```

end
theory Complex-Axial-Symmetry

imports Complex-Angles Complex-Triangles

begin

```

6 Axial symmetry in complex field

In the following we define the axial symmetry and prove basics properties.

```

context
  fixes z1 z2 :: complex and  $\alpha \ \beta :: \text{complex}$ 
  assumes neq0:  $\langle z1 \neq z2 \rangle$ 
  defines  $\langle \alpha \equiv (z1 - z2) / (\text{cnj } z1 - \text{cnj } z2) \rangle$ 
  defines  $\langle \beta \equiv (z2 * \text{cnj } z1 - z1 * \text{cnj } z2) / (\text{cnj } z1 - \text{cnj } z2) \rangle$ 
begin

definition axial-symmetry:  $\langle \text{complex} \Rightarrow \text{complex} \rangle$  where
   $\langle \text{axial-symmetry } z \equiv \text{cnj } z * (z1 - z2) / (\text{cnj } z1 - \text{cnj } z2) + (z2 * \text{cnj } z1 - z1 * \text{cnj } z2) / (\text{cnj } z1 - \text{cnj } z2) \rangle$ 

lemma norm- $\alpha$ -eq-1:  $\langle \text{cmod } (\alpha) = 1 \rangle$ 
  by (auto simp: neq0  $\alpha$ -def  $\beta$ -def)

lemma z1-inv:  $\langle \text{axial-symmetry } z1 = z1 \rangle$ 
  unfolding axial-symmetry-def
  by (metis (mono-tags, lifting) add-diff-eq add-divide-distrib complex-cnj-cancel-iff

    diff-add-cancel eq-iff-diff-eq-0 mult.commute neq0 nonzero-eq-divide-eq right-diff-distrib')

lemma z2-inv:  $\langle \text{axial-symmetry } z2 = z2 \rangle$ 

```

unfolding *axial-symmetry-def*
by (*smt* (*z3*) *add-diff-cancel-left add-diff-cancel-left' add-divide-distrib axial-symmetry-def*
complex-cnj-cancel-iff diff-add-cancel divide-divide-eq-right eq-divide-imp mult-commute-abs
right-minus-eq z1-inv)

lemma *cmod-axial*: $\langle \text{cmod } (\text{axial-symmetry } z - \text{axial-symmetry } z') = \text{cmod } (\alpha * (\text{cnj } z - \text{cnj } z')) \rangle$
unfolding *axial-symmetry-def*
by (*auto simp: α -def β -def*) (*simp add: diff-divide-distrib mult.commute right-diff-distrib'*)

lemma *cmod-axial-inv*: $\langle \text{cmod } (\text{axial-symmetry } z - \text{axial-symmetry } z') = \text{cmod } (z - z') \rangle$
by (*metis cmod-axial complex-cnj-diff complex-mod-cnj mult.commute*
mult.right-neutral norm- α -eq-1 norm-mult)

lemma *axial-symmetry-dist1*: $\langle \text{cdist } z1 \ z = \text{cdist } z1 \ (\text{axial-symmetry } z) \rangle$
by (*metis cdist-def cmod-axial-inv z1-inv*)

lemma *axial-symmetry-dist2*: $\langle \text{cdist } z2 \ z = \text{dist } z2 \ (\text{axial-symmetry } z) \rangle$
by (*metis cdist-def cmod-axial-inv dist-commute dist-norm z2-inv*)

lemma $\alpha\beta$: $\langle \alpha * \text{cnj } \beta + \beta = 0 \rangle$
by (*simp add: add-divide-eq-iff α -def β -def*)

lemma *involution-symmetry*: $\langle \text{axial-symmetry } (\text{axial-symmetry } z) = z \rangle$
proof –
have $\langle \alpha * \text{cnj } (\alpha * \text{cnj } z + \beta) + \beta = \alpha * \text{cnj } \alpha * z + \alpha * \text{cnj } \beta + \beta \rangle$
by (*simp add: ring-class.ring-distrib(1)*)
then show ?thesis **unfolding** *axial-symmetry-def*
by (*smt (verit, best) $\alpha\beta$ add.right-neutral cmod-eq-one group-cancel.add1*
mult.commute mult-cancel-right2 norm- α -eq-1 times-divide-eq-left α -def
 β -def)
qed

lemma *arg- α* : $\langle \text{Arg } \alpha = \lfloor 2 * \text{Arg } (z1 - z2) \rfloor \rangle$
unfolding α -def β -def
by (*smt (verit, del-insts) arg-cnj-not-pi arg-div arg-pi-iff canon-ang-id canon-ang-pi-3pi*
complex-cnj-cancel-iff complex-cnj-diff eq-cnj-iff-real eq-iff-diff-eq-0 neq0 pi-ge-two)

lemma *Arg-invol*: $\langle \text{Arg } (\text{axial-symmetry } (\text{axial-symmetry } z)) = \text{Arg } z \rangle$
by (*simp add: involution-symmetry*)

lemma *angle-sum-symmetry*: $\langle z \neq z1 \implies \lfloor \text{angle-c } z \ z1 \ z2 + \text{angle-c } z2 \ z1 \ (\text{axial-symmetry } z) \rfloor = \text{angle-c } z \ z1 \ (\text{axial-symmetry } z) \rangle$
proof –
assume $\langle z \neq z1 \rangle$

have $\langle z2 - z1 \neq 0 \rangle$ **using** *angle-c-sum neq0* **by** *auto*
have $\langle z1 - (\text{axial-symmetry } z) \neq 0 \rangle$
by (*metis* $\langle z \neq z1 \rangle$ *eq-iff-diff-eq-0 involution-symmetry z1-inv*)
then show *?thesis*
using *angle-c-sum*
by (*metis* $\langle z2 - z1 \neq 0 \rangle$ *right-minus-eq z1-inv*)
qed

lemma *angle-symmetry-eq-imp*:
assumes $h: \langle z1 \neq z \rangle \langle z2 \neq z \rangle$
shows $\langle \text{angle-c } z \ z1 \ z2 = - \text{angle-c } (\text{axial-symmetry } z) \ z1 \ z2 \vee \text{angle-c } z \ z1 \ z2 = \text{angle-c } (\text{axial-symmetry } z) \ z1 \ z2 \rangle$
by (*metis* (*mono-tags, lifting*) *axial-symmetry-dist1 cdist-def cmod-axial-inv congruent-ctriangle-sss(22)* $h(1)$ *neq0 z2-inv*)

lemma *angle-symmetry*:
assumes $h: \langle z1 \neq z \rangle \langle z2 \neq z \rangle$
and $\langle \text{angle-c } z \ z1 \ z2 = \text{angle-c } (\text{axial-symmetry } z) \ z1 \ z2 \rangle$
shows $\langle z = \text{axial-symmetry } z \rangle$
proof –
have $\langle \text{cmod } (z - z1) = \text{cmod } (\text{axial-symmetry } z - z1) \rangle$
using *axial-symmetry-dist1* **by** *auto*
have $\langle \text{cmod } (z - z2) = \text{cmod } (\text{axial-symmetry } z - z2) \rangle$
by (*metis* *axial-symmetry-dist2 cdist-def dist-commute dist-complex-def*)
have $\langle z - z1 = \text{axial-symmetry } z - z1 \rangle$
by (*metis* $\langle \text{cmod } (z - z1) = \text{cmod } (\text{local.axial-symmetry } z - z1) \rangle$ *ang-cos ang-sin asms(3)* *complex-cnj-cnj complex-eq-iff eq-iff-diff-eq-0 mult-cancel-left neq0*)
then show *?thesis*
by *simp*
qed

lemma *line-is-inv*: $\langle z \in \text{line } z1 \ z2 \wedge z \neq z2 \wedge z \neq z1 \implies z = \text{axial-symmetry } z \rangle$
proof –
assume $\langle z \in \text{line } z1 \ z2 \wedge z \neq z2 \wedge z \neq z1 \rangle$
then have $\langle \text{angle-c } z \ z1 \ z2 = 0 \vee \text{angle-c } z \ z1 \ z2 = \pi \rangle$
unfolding *line-def* **using** *neq0 collinear-angle*
using *collinear-sym1 collinear-sym2'* **by** *blast*
then show *?thesis*
by (*smt* (*verit*) $\langle z \in \text{line } z1 \ z2 \wedge z \neq z2 \wedge z \neq z1 \rangle$ *angle-c-commute angle-c-commute-pi angle-symmetry angle-symmetry-eq-imp*)
qed

lemma *dist-inv*: $\langle \text{cdist } a \ b = \text{cdist } (\text{axial-symmetry } a) \ (\text{axial-symmetry } b) \rangle$
by (*simp* *add: cmod-axial-inv*)

lemma *collinear-inv*: **assumes** $\langle \text{collinear } a \ b \ c \rangle$ **and** $\langle a \neq b \wedge b \neq c \wedge c \neq a \rangle$
shows $\langle \text{collinear } (\text{axial-symmetry } a) \ (\text{axial-symmetry } b) \ (\text{axial-symmetry } c) \rangle$
proof –

have $\langle \text{angle-}c\ a\ b\ c = \pi \vee \text{angle-}c\ a\ b\ c = 0 \rangle$
using *assms(1) assms(2) collinear-angle* **by** *blast*
then have $\langle \text{angle-}c\ (\text{axial-symmetry}\ a)\ (\text{axial-symmetry}\ b)\ (\text{axial-symmetry}\ c)$
 $= \text{angle-}c\ a\ b\ c$
 $\vee \text{angle-}c\ (\text{axial-symmetry}\ a)\ (\text{axial-symmetry}\ b)\ (\text{axial-symmetry}\ c) = -$
 $\text{angle-}c\ a\ b\ c \rangle$
by (*metis congruent-ctriangle-sss(24) assms(2) dist-inv minus-equation-iff*)
then show *?thesis*
using $\langle \text{angle-}c\ a\ b\ c = \pi \vee \text{angle-}c\ a\ b\ c = 0 \rangle$ *collinear-sin-neq-0 collinear-sym1*
by *fastforce*
qed

lemma *axial-symmetry-eq-line*: $\langle z \neq z1 \wedge z \neq z2 \implies z = \text{axial-symmetry}\ z \implies z \in$
 $\text{line}\ z1\ z2 \rangle$

proof $-$

assume $\langle z \neq z1 \wedge z \neq z2 \rangle$ $\langle z = \text{axial-symmetry}\ z \rangle$
then have $g0: \langle z = \alpha * \text{cnj}\ z + \beta \rangle$ **unfolding** *axial-symmetry-def*
by (*simp add: mult.commute α -def β -def*)
then have $g1: \langle \text{cnj}\ z = (z - \beta) * \text{cnj}\ \alpha \rangle$
by (*smt (verit, ccfv-threshold) add-diff-cancel complex-cnj-cnj complex-cnj-diff*
 α -def β -def
 $\text{complex-cnj-divide mult.commute neq0 nonzero-eq-divide-eq right-minus-eq}$
 $\text{times-divide-eq-left}$)
also have $g2: \langle z = (\text{cnj}\ z - \text{cnj}\ \beta) * \alpha \rangle$
by (*metis calculation complex-cnj-cnj complex-cnj-diff complex-cnj-mult*)
have $\langle \alpha / (z1 - z2) = 1 / (\text{cnj}\ z1 - \text{cnj}\ z2) \rangle$
by (*auto simp: α -def β -def*)
have $\langle \text{Im}\ w = (w - \text{cnj}\ w) / (2 * i) \rangle$ **for** w
using *Im-express-cnj* **by** *blast*
then have $\langle \text{Im}\ (z2 * \text{cnj}\ z1) = (z2 * \text{cnj}\ z1 - \text{cnj}\ (z2 * \text{cnj}\ z1)) / (2 * i) \rangle$
by *presburger*
then have $\langle \text{Im}\ (z2 * \text{cnj}\ z1) * 2 * i = (z2 * \text{cnj}\ z1 - \text{cnj}\ z2 * z1) \rangle$
by (*metis complex-cnj-cnj complex-cnj-mult complex-diff-cnj mult.commute*)
have $f0: \langle (\text{cnj}\ z1 - \text{cnj}\ z2) * (z1 - z2) = \text{cmod}\ z1^2 + \text{cmod}\ z2^2 - \text{cnj}\ z1 * z2$
 $- \text{cnj}\ z2 * z1 \rangle$
by (*smt (verit) cancel-ab-semigroup-add-class.diff-right-commute complex-mult-cnj-cmod*

$\text{diff-diff-eq2 left-diff-distrib mult.assoc mult.commute norm-minus-commute}$
 $\text{norm-mult of-real-add of-real-eq-iff}$)

have $f1: \langle \text{cmod}\ z1^2 + \text{cmod}\ z2^2 - \text{cnj}\ z1 * z2 - \text{cnj}\ z2 * z1 = \text{cmod}\ z1^2 +$
 $\text{cmod}\ z2^2 - 2 * \text{Re}\ z1 * \text{Re}\ z2 - 2 * \text{Im}\ z1 * \text{Im}\ z2 \rangle$

by (*auto simp: field-simps intro: complex-eqI*)

have $\langle \beta = 2 * i * \text{Im}\ (z2 * \text{cnj}\ z1) / (\text{cnj}\ z1 - \text{cnj}\ z2) \rangle$

using $\langle \text{cor}\ (\text{Im}\ (z2 * \text{cnj}\ z1)) = (z2 * \text{cnj}\ z1 - \text{cnj}\ (z2 * \text{cnj}\ z1)) / (2 * i) \rangle$

unfolding α -def β -def **by** *auto*

then have $f2: \langle \beta / (z1 - z2) = 2 * i * \text{Im}\ (z2 * \text{cnj}\ z1) / ((\text{cnj}\ z1 - \text{cnj}\ z2) * (z1 - z2)) \rangle$

by *simp*

then have $\langle \beta / (z1 - z2) = 2 * i * \text{Im}\ (z2 * \text{cnj}\ z1) / (\text{cmod}\ z1^2 + \text{cmod}\ z2^2 -$
 $2 * \text{Re}\ z1 * \text{Re}\ z2 - 2 * \text{Im}\ z1 * \text{Im}\ z2) \rangle$

```

    using f0 f1 by presburger
  then have ⟨is-real ((z - z1)/(z1-z2))⟩
    unfolding axial-symmetry-def using g2
  by (smt (verit, del-insts) add-diff-cancel-right axial-symmetry-def complex-cnj-cnj
      complex-cnj-diff complex-cnj-divide diff-divide-distrib divide-divide-eq-right
eq-cnj-iff-real
      g0 g2 mult.commute times-divide-eq-right z1-inv z2-inv α-def β-def)
  then show ⟨z ∈ line z1 z2⟩
    unfolding line-def collinear-def
  by (metis (mono-tags, lifting) Im-i-times Re-i-times cnj.simps(2) complex-i-mult-minus
      eq-cnj-iff-real mem-Collect-eq minus-diff-eq minus-divide-right)
qed

lemma angle-symmetry-eq:
  assumes h:⟨z1 ≠ z⟩ ⟨z2 ≠ z⟩ ⟨z ∉ line z1 z2⟩
  shows ⟨angle-c z z1 z2 = - angle-c (axial-symmetry z) z1 z2⟩
proof -
  have f0:⟨angle-c z z1 z2 = - angle-c (axial-symmetry z) z1 z2 ∨
      angle-c z z1 z2 = angle-c (axial-symmetry z) z1 z2⟩
    using angle-symmetry-eq-imp h(1) h(2) by blast
  have f1:⟨angle-c z z2 z1 = - angle-c (axial-symmetry z) z2 z1
      ∨ angle-c z z2 z1 = angle-c (axial-symmetry z) z2 z1⟩
    by (metis axial-symmetry-dist1 congruent-ctriangle-sss(24) dist-inv h(1) neq0
z2-inv)
  show ?thesis
    using angle-symmetry axial-symmetry-eq-line f0 h(1) h(2) h(3) by presburger
qed

end
end
theory Morley

```

```

imports Complex-Axial-Symmetry

```

```

begin

```

7 Rotations

```

locale complex-rotation =
  fixes A::complex and ϑ::real
begin

```

```

definition ⟨r z = A + (z-A)*cis(ϑ)⟩

```

```

lemma cmod-inv-rotation:⟨cmod (z-A) = cmod (r z - A)⟩
  unfolding r-def
  by (simp add: norm-mult)

```



```

lemma inner-ang: $\langle \cos (\angle z1\ z2) * (cmod\ z1\ * cmod\ z2) = Re\ (innerprod\ z1\ z2) \rangle$ 
proof -
  have  $\langle Re\ (innerprod\ z1\ z2) = Re\ (scalprod\ z1\ z2) \rangle$ 
    unfolding innerprod-def by (auto)
  then show ?thesis
    by (metis cos-cmod-scalprod mult.commute)
qed

lemma ang-eq-cos-theta: $\langle z \neq A \implies \cos (angle-c\ z\ A\ (r\ z)) = \cos (\vartheta) \rangle$ 
proof -
  assume  $\langle z \neq A \rangle$ 
  then have  $\langle innerprod\ (z-A)\ ((r\ z - A)) = (z-A) * cis(\vartheta) * cnj(z - A) \rangle$ 
    unfolding innerprod-def r-def by auto
  then have  $f0: \langle innerprod\ (z-A)\ ((r\ z - A)) = cmod\ (z-A)^2 * cis(\vartheta) \rangle$ 
    by (metis ab-semigroup-mult-class.mult-ac(1) complex-mult-cnj-cmod mult.commute)
  then have  $\langle \cos (angle-c\ z\ A\ (r\ z)) * cmod\ (z-A) * cmod(z - A) = Re\ (innerprod\ (z-A)\ ((r\ z - A))) \rangle$ 
    unfolding angle-c-def
    by (metis inner-ang mult.assoc cmod-inv-rotation)
  then have  $\langle \cos (angle-c\ z\ A\ (r\ z)) = Re\ (cis\ \vartheta) \rangle$ 
    by (simp add:  $\langle z \neq A \rangle$  f0 power2-eq-square)
  then show ?thesis
    by (auto simp: cis.code)
qed

lemma cdist-dist: $\langle cdist = dist \rangle$ 
  using cdist-commute dist-complex-def by fastforce

lemma ang-eq-theta: assumes  $h: \langle z \neq A \rangle$  shows  $\langle angle-c\ z\ A\ (r\ z) = |\vartheta| \rangle$ 
proof (cases  $\langle angle-c\ z\ A\ (r\ z) = |\vartheta| \rangle$ )
  case True
    then have  $\langle r\ z = A + (z-A) * cis(-\vartheta) \rangle$ 
      by (metis add-diff-cancel-left' arg-cis cdist-def cis-cmod cis-neq-zero cmod-inv-rotation)
    
$$isoscele-iff-pr-cis-qr\ mult.left-commute\ mult.right-neutral\ mult.cancel-left$$

    
$$norm-cis$$

    
$$norm-minus-commute\ of-real-1\ r-def\ right-minus-eq$$

    then show ?thesis
      by (smt (verit, del-insts) True add-diff-cancel-left' angle-c-neq0 arg-cis arg-mult-eq
        canon-ang-cos canon-ang-sin cis.code cis-cnj diff-add-cancel divisors-zero
r-def)
  next
    case False
    have  $\langle \cos (angle-c\ z\ A\ (r\ z)) = \cos (\vartheta) \rangle$ 
      using ang-eq-cos-theta h by auto
    then show ?thesis
      by (smt (verit, ccfv-threshold) cmod-inv-rotation r-def add-diff-cancel-left' an-
gle-c-neq0)

```

$\text{angle-c-sum } \text{arg-cis } \text{canon-ang-sin } \text{cdist-def } \text{cis.code } \text{cis-neq-zero } \text{divisors-zero}$
 $\text{eq-iff-diff-eq-0 } h$
 $\text{isoscele-iff-pr-cis-qr } \text{nonzero-mult-div-cancel-left } \text{nonzero-mult-div-cancel-right}$
 $\text{norm-minus-commute})$
qed

lemma $\text{inj-r} : \langle \text{inj } r \rangle$
unfolding inj-on-def **by** (auto simp:r-def)

lemma $\text{img-eqI} : \langle \text{cdist } A \ z1 = \text{cdist } A \ z2 \wedge \text{angle-c } z1 \ A \ z2 = \vartheta \implies z2 = r \ z1 \rangle$
apply $(\text{cases } \langle z1 = A \vee z2 = A \rangle)$
using $r\text{-def } \text{add-minus-cancel } \text{isoscele-iff-pr-cis-qr}$ **apply** force
unfolding $r\text{-def}$
by $(\text{metis } \text{add.commute } \text{cdist-commute } \text{diff-add-cancel } \text{isoscele-iff-pr-cis-qr } \text{mult.commute})$

lemma $\text{r-id-iff} : \langle |\vartheta| = 0 \longleftrightarrow r = \text{id} \rangle$
proof –
obtain $cc :: (\text{complex} \Rightarrow \text{complex}) \Rightarrow \text{complex}$
and $cca :: (\text{complex} \Rightarrow \text{complex}) \Rightarrow \text{complex}$ **where**
 $f1 : \forall f. (\text{id} = f \vee f \ (cc \ f) \neq cc \ f) \wedge ((\forall c. f \ c = c) \vee \text{id} \neq f)$
by $(\text{metis } (\text{no-types}) \ \text{eq-id-iff})$
have $f2 : \forall c \ ca \ ra. ca + (c - ca) * \text{Complex } (\cos \ ra) \ (\sin \ ra) = \text{complex-rotation.r}$
 $ca \ ra \ c$
by $(\text{simp add: complex-rotation.r-def cis.code})$
then have $\forall c \ ca. \text{complex-rotation.r } c \ 0 \ ca = ca$
by $(\text{metis } (\text{no-types}) \ \text{add-diff-cancel-left' } \cos\text{-zero } \text{diff-diff-eq2}$
 $\lambda\text{mbda-one } \text{mult.commute } \text{one-complex.code } \sin\text{-zero})$
then show $?thesis$
using $f2 \ f1$
by $(\text{metis } (\text{lifting, full-types}) \ \text{ang-eq-theta } \text{angle-c-neq0}$
 $\text{canon-ang-cos } \text{canon-ang-sin } \text{complex-i-not-zero})$
qed
end

lemma $\text{axial-symmetry-eq} : \langle \text{axial-symmetry } B \ C \ P = \text{axial-symmetry } C \ B \ P \rangle$ **if**
 $\langle C \neq B \rangle$ **for** $C \ B \ P$
unfolding $\text{axial-symmetry-def}[OF \ \text{that}] \ \text{axial-symmetry-def}[OF \ \text{that}[\text{symmetric}]]$
by $(\text{metis } (\text{no-types, lifting}) \ \text{complex-cnj-cancel-iff } \text{eq-iff-diff-eq-0}$
 $\text{minus-diff-eq } \text{nonzero-minus-divide-divide } \text{times-divide-eq-right})$

lemma $\text{img-r-sym} :$
assumes $h : \langle z1 \neq z2 \rangle \ \langle z \notin \text{line } z1 \ z2 \rangle$
shows $\langle \text{axial-symmetry } z1 \ z2 \ z = \text{complex-rotation.r } z1 \ (\lfloor 2 * \text{angle-c } z \ z1 \ z2 \rfloor) \ z \rangle$
proof –
interpret $\text{complex-rotation } z1 \ \lfloor 2 * \text{angle-c } z \ z1 \ z2 \rfloor .$
let $?z = \langle \text{axial-symmetry } z1 \ z2 \ z \rangle$

```

from h have ⟨z1 ≠ z⟩ ⟨z2 ≠ z⟩
  unfolding line-def Elementary-Complex-Geometry.collinear-def by(auto)
then have ⟨angle-c ?z z1 z2 = -angle-c z z1 z2⟩
  using angle-symmetry-eq h(1) h(2) by force
then have ⟨angle-c z z1 ?z = |2*angle-c z z1 z2|⟩
  by (metis ⟨z1 ≠ z⟩ add.inverse-inverse angle-c-commute angle-c-commute-pi
    angle-sum-symmetry h(1) mult-2-right mult-commute-abs)
have ⟨cdist z1 z = cdist z1 (axial-symmetry z1 z2 z)⟩
  using axial-symmetry-dist1 h(1) by blast
then show ?thesis
  apply(intro img-eqI)
  by (metis ⟨angle-c z z1 (axial-symmetry z1 z2 z) = |2 * angle-c z z1 z2|⟩)
qed

```

```

lemma img-r-sym':
  assumes h:⟨z1 ≠ z2⟩ ⟨z ∉ line z1 z2⟩
  shows ⟨axial-symmetry z1 z2 z = complex-rotation.r z1 (|-2*angle-c z2 z1 z|)
    z⟩
  by (metis angle-c-commute angle-c-neq0 axial-symmetry-eq-line
    cancel-comm-monoid-add-class.diff-cancel complex-rotation.img-eqI
    complex-rotation.r-def h(1,2) img-r-sym mult-eq-0-iff mult-minus-left
    mult-minus-right pi-neq-zero two-pi-canonical)

```

```

lemma equality-for-pqr:
  assumes 1:⟨(a2::complex)*a3 ≠ 1⟩ and 2:⟨∧z. h z = a3*z + b3⟩ and 3:⟨∧z. g
    z = a2*z + b2⟩ and 4:⟨g (h z) = z⟩
  shows ⟨z = (a2*b3 + b2)/(1-a2*a3)⟩
proof -
  have f21:⟨g (h z) = a2*a3*z + a2*b3 + b2⟩
    using assms by(auto simp:2 3field-simps)
  then have ⟨g (h z) = a2*a3*z + a2*b3 + b2 ⟷ z*(1-a2*a3) = a2*b3 +
    b2⟩
  by(auto simp:field-simps 4)
  then have ⟨a2*a3 ≠ 1 ⟹ z = (a2*b3 + b2)/(1-a2*a3)⟩
    using f21 by(auto simp:field-simps)
  then show ?thesis using 1 by auto
qed

```

```

lemma equality-for-comp:
  assumes 2:⟨∧z. h z = (a3::complex)*z + b3⟩ and 3:⟨∧z. g z = a2*z + b2⟩
  and 4:⟨∧z. f z = a1*z + b1⟩
  shows ⟨((f∘f∘f)∘(g∘g∘g)∘(h∘h∘h)) z = (a1*a2*a3)^3*z + (a1^2+a1+1)*b1
    + a1^3*(a2^2+a2+1)*b2
    + a1^3*a2^3*(a3^2+a3+1)*b3 ⟩
  using assms unfolding comp-def by(auto simp:fun-eq-iff power2-eq-square power3-eq-cube
    field-simps)

```

```

lemma eq-translation-id:
  assumes ⟨h = complex-rotation.r A 0⟩ ⟨h B = B⟩

```

```

shows ⟨h = id⟩
using assms(1) complex-rotation.r-id-iff by auto

lemma r-eqI:
  assumes ⟨A = B⟩ ⟨∅1 = ∅2⟩
  shows ⟨r A ∅1 = r B ∅2⟩
  using assms(1) assms(2) by force

lemma r-eqI':
  assumes ⟨A = B⟩ ⟨∅1 = ∅2⟩
  shows ⟨r A ∅1 z = r B ∅2 z⟩
  using assms(1) assms(2) by force

lemma composed-rotations-same-center:
  shows ⟨(complex-rotation.r A ∅1 ∘ complex-rotation.r A ∅2) = complex-rotation.r
  A (∅1 + ∅2)⟩
  unfolding complex-rotation.r-def by (auto simp: fun-eq-iff cis-mult add-ac)

lemma composed-rotations:
  assumes h:⟨|∅1 + ∅2| ≠ 0⟩
  shows ⟨(complex-rotation.r A ∅1 ∘ complex-rotation.r B ∅2) =
  complex-rotation.r ((A*(1-cis ∅1) + B*cis ∅1*(1-cis ∅2))/(1-cis
  (∅1+∅2))) (∅1 + ∅2)⟩
  proof -
    have ⟨cis (∅1 + ∅2) ≠ 1⟩
      by (metis arg-cis assms cis-zero zero-canonical)
    with h have ⟨(complex-rotation.r A ∅1 ∘ complex-rotation.r B ∅2)
      ((A*(1-cis ∅1) + B*cis ∅1*(1-cis ∅2))/(1-cis (∅1+∅2)))
      = (A*(1-cis ∅1) + B*cis ∅1*(1-cis ∅2))/(1-cis (∅1+∅2))⟩
      unfolding complex-rotation.r-def using assms
      by (auto simp: cis-mult field-simps intro!:)
    with h show ?thesis
      apply (cases ⟨∅1 = 0 ∨ ∅2 = 0⟩)
      unfolding complex-rotation.r-def using assms
      by (auto simp: cis-mult field-simps fun-eq-iff diff-divide-distrib add-divide-distrib
      intro!:)
  qed

lemma composed-rotation-is-trans:
  assumes ⟨|∅1 + ∅2| = 0⟩
  shows ⟨(complex-rotation.r A ∅1 ∘ complex-rotation.r B ∅2) z = z + (B -
  A)*(cis(∅1) - 1)⟩
  using assms
  by (auto simp: complex-rotation.r-def add-divide-distrib diff-divide-distrib field-simps)

  (metis add.commute canon-ang-cos canon-ang-sin cis.code cis-mult cis-zero
  lambda-one mult.commute)

```

8 Morley's theorem

We begin by proving the Morley's theorem in the case where angles are positives then using the congruence between two triangles with the same angles only not of the same sign we prove Morley's theorem when angles are negatives.

We then proceed to conclude because in a triangle either angles are all negatives or all the angles are positives depending on orientation.

theorem *Morley-pos:*

```

assumes  $\neg \text{collinear } A \ B \ C$ 
   $\langle \text{angle-c } A \ B \ R = \text{angle-c } A \ B \ C / 3 \rangle$  (is  $\langle ?abr = ?abc \rangle$ )
   $\langle \text{angle-c } B \ A \ R = \text{angle-c } B \ A \ C / 3 \rangle$  (is  $\langle ?bar = ?\alpha \rangle$ )
   $\langle \text{angle-c } B \ C \ P = \text{angle-c } B \ C \ A / 3 \rangle$  (is  $\langle ?bcp = ?bca \rangle$ )
   $\langle \text{angle-c } C \ B \ P = \text{angle-c } C \ B \ A / 3 \rangle$  (is  $\langle ?cbp = ?\beta \rangle$ )
   $\langle \text{angle-c } C \ A \ Q = \text{angle-c } C \ A \ B / 3 \rangle$  (is  $\langle ?caq = ?cab \rangle$ )
   $\langle \text{angle-c } A \ C \ Q = \text{angle-c } A \ C \ B / 3 \rangle$  (is  $\langle ?acq = ?\gamma \rangle$ )
  and  $hhh: \langle \text{angle-c } B \ A \ C / 3 + \text{angle-c } C \ B \ A / 3 + \text{angle-c } A \ C \ B / 3 \rangle = \text{pi} / 3$ 
shows  $\langle \text{cdist } R \ P = \text{cdist } P \ Q \wedge \text{cdist } Q \ R = \text{cdist } P \ Q \rangle$ 
proof -
  have  $\text{bundle-line}: \langle A \notin \text{line } B \ C \rangle \langle B \notin \text{line } A \ C \rangle \langle C \notin \text{line } A \ B \rangle \langle A \neq B \rangle \langle B \neq C \rangle$ 
   $\langle C \neq A \rangle$ 
  using  $\text{assms}(1)$  non-collinear-independant by (auto simp: collinear-sym1 collinear-sym2 line-def)
  {fix  $A \ B \ C \ \gamma$ 
    assume  $ABC\text{-nline}: \langle A \notin \text{line } B \ C \rangle$ 
    and  $eq\text{-}3c: \langle \text{angle-c } A \ C \ B = 3 * \gamma \rangle$ 
    then have  $neg\text{-}PI: \langle \text{abs } \gamma < \text{pi} / 3 \rangle$ 
    proof -
      have  $\langle \text{angle-c } A \ C \ B \neq \text{pi} \rangle$ 
      using  $ABC\text{-nline}(1)$  unfolding line-def
      by (metis angle-c-commute-pi collinear-iff mem-Collect-eq non-collinear-independant)
      then have  $\langle \text{angle-c } A \ C \ B \in \{-\text{pi} < .. < \text{pi}\} \rangle$ 
      using ang-c-in less-eq-real-def by auto
      then show  $\langle \text{abs } \gamma < \text{pi} / 3 \rangle$ 
      using eq-3c by force
    qed} note  $ang\text{-}inf\text{-}pi3 = \text{this}$ 
  have  $\langle \text{angle-c } B \ A \ C + \text{angle-c } C \ B \ A + \text{angle-c } A \ C \ B \rangle = \text{pi}$ 
  by (metis collinear-def add.commute angle-sum-triangle-c assms(1) collinear-sym1 collinear-sym2)
  have  $eq\text{-}pi: \langle 3 * ?\alpha + 3 * ?\beta + 3 * ?\gamma \rangle = \text{pi}$ 
  using  $\langle \text{angle-c } B \ A \ C + \text{angle-c } C \ B \ A + \text{angle-c } A \ C \ B \rangle = \text{pi}$  by force
  then have  $neg\text{-}pi: \langle \text{abs } ?\beta \neq \text{pi} \wedge \text{abs } ?\gamma \neq \text{pi} \wedge \text{abs } ?\alpha \neq \text{pi} \rangle$ 
  by (smt (verit) ang-neg-neg angle-c-commute angle-c-neq0 assms canon-ang-sin collinear-angle collinear-sin-neq-0 divide-eq-0-iff divide-nonneg-pos)

   $\text{minus-divide-left } \text{pi-neq-zero } \text{sin-pi } \text{sin-pi-minus zero-canonical}$ 
  moreover have  $\langle 3 * ?\alpha \neq 0 \wedge 3 * ?\beta \neq 0 \wedge 3 * ?\gamma \neq 0 \rangle$ 
  using bundle-line collinear-sin-neq-0 line-def angle-c-commute assms(1) bun-

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dle-line(4)
  bundle-line(5) bundle-line(6) collinear-iff assms(1) collinear-sin-neq-0
  by (metis collinear-sym2' divide-eq-eq-numeral1(1) mem-Collect-eq mult.commute
zero-neq-numeral)
  ultimately have neq-0:⟨ $\lfloor ?\beta \rfloor \neq 0 \wedge \lfloor ?\gamma \rfloor \neq 0 \wedge \lfloor ?\alpha \rfloor \neq 0$ ⟩
  by (metis ang-vec-def angle-c-def arg-cis assms(3) assms(5) assms(7) canon-ang-arg
mult-zero-right)

interpret rot1: complex-rotation A 2*?α .
interpret rot2: complex-rotation B 2*?β .
interpret rot3: complex-rotation C 2*?γ .

let ?f=⟨rot1.r⟩
have f0:⟨rot1.r A = A⟩
  unfolding rot1.r-def by auto
have f1:⟨rot2.r B = B⟩
  unfolding rot2.r-def by auto
have f2:⟨rot3.r C = C⟩
  unfolding rot3.r-def by auto

have ⟨cmod (rot3.r z - C) = cmod (z-C)⟩ for z
  using rot3.cmod-inv-rotation by presburger
have f2:⟨B≠C⟩
  by (metis collinear-def assms(1) collinear-sym1 collinear-sym2)
have f5:⟨angle-c C B P = ?β⟩ ⟨angle-c P B C = -?β⟩
  by (auto simp add: assms(5) angle-c-commute)
  (metis angle-c-commute angle-c-commute-pi assms(5) neq-pi pi-canonical)
then have f3:⟨P ∉ line C B⟩
  by (smt (verit, ccfv-SIG) neq-0 neq-pi angle-c-commute angle-c-neq0 assms(4)
collinear-angle divide-eq-0-iff line-def mem-Collect-eq pi-canonical zero-canonical)
then have f3':⟨P ∉ line B C⟩
  using collinear-sym2' line-def by blast
then have f4:⟨P≠C ∧ P≠B⟩
  by (metis collinear-def collinear-sym1 collinear-sym2 line-def mem-Collect-eq)
then have ⟨angle-c P B C = - angle-c (axial-symmetry B C P) B C⟩
  using angle-symmetry-eq[OF f2 - f3'] by auto
then have ⟨angle-c (axial-symmetry B C P) B C = ?β⟩
  using f5(2) by fastforce

have f13:⟨angle-c P C B = ?γ⟩
  by (smt (verit, ccfv-threshold) angle-c-commute angle-c-commute-pi assms(1)
assms(4) assms(7)
collinear-sin-neq-0 nonzero-minus-divide-divide nonzero-minus-divide-right
sin-pi)
let ?P'=⟨(axial-symmetry B C P)⟩

have ⟨angle-c (axial-symmetry B C P) B P =  $\lfloor 2*?\beta \rfloor$ ⟩
  by (metis ⟨P ≠ C ∧ P ≠ B⟩ ⟨angle-c (axial-symmetry B C P) B C = angle-c
C B A / 3⟩)

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    angle-sum-symmetry f2 f5(1) involution-symmetry mult.commute mult-2-right
z1-inv)
  have ⟨cdist B (P) = cdist B ?P'⟩
  by (auto)(metis cmod-axial-inv f2 z1-inv)
  then have ⟨cdist B (rot2.r P) = cdist B ?P'⟩
  unfolding cdist-def
  using rot2.cmod-inv-rotation by presburger
  have f16:⟨rot2.r ?P' = P⟩
  by (metis ⟨angle-c (axial-symmetry B C P) B P = |2 * (angle-c C B A / 3)|⟩
    ⟨cdist B P = cdist B (axial-symmetry B C P)⟩ canon-ang-cos
    canon-ang-sin cis.code complex-rotation.img-eqI complex-rotation.r-def)
  have f10:⟨angle-c P C B = |?γ|⟩
  by (metis ⟨angle-c P C B = angle-c A C B / 3⟩ ang-vec-def angle-c-def arg-cis
    canon-ang-arg)

  from f10 have f11:⟨angle-c B C ?P' = |?γ|⟩
  by (metis axial-symmetry-eq ⟨P ≠ C ∧ P ≠ B⟩ ⟨angle-c P C B = angle-c A C
    B / 3⟩ angle-c-commute
    angle-symmetry angle-symmetry-eq-imp axial-symmetry-eq-line canon-ang-uminus-pi
    f2 f3 pi-canonical)
  then have f12:⟨angle-c P C ?P' = |2*?γ|⟩
  by (metis axial-symmetry-eq f13 angle-sum-symmetry f10 f2 f4 mult.commute
    mult-2-right)
  then have f15:⟨rot3.r P = ?P'⟩
  by (metis axial-symmetry-eq canon-ang-cos canon-ang-sin cis.code f13 f2 f3
    img-r-sym complex-rotation.r-def)
  have P-inv:⟨rot2.r (rot3.r P) = P⟩
  using f16 f15 by presburger
  let ?Q' = ⟨axial-symmetry A C Q⟩
  have ⟨angle-c C A Q = -?α⟩
  by (metis angle-c-commute assms(1) assms(6) collinear-sin-neq-0
    collinear-sym1 collinear-sym2' minus-divide-left sin-pi)
  then have ⟨angle-c Q A C = ?α⟩
  by (metis add.inverse-inverse angle-c-commute canon-ang-uminus-pi neq-pi
    pi-canonical)
  have ⟨A ≠ C ∧ Q ∉ line A C⟩
  by (metis ⟨angle-c Q A C = angle-c B A C / 3⟩ angle-c-neq0 assms(7)
    collinear-angle div-0
    f10 f13 line-def mem-Collect-eq neq-0 neq-pi zero-canonical)
  then have f17:⟨rot1.r Q = ?Q'⟩
  using img-r-sym[of A C Q]
  using ⟨angle-c Q A C = angle-c B A C / 3⟩ canon-ang-cos canon-ang-sin
    cis.code complex-rotation.r-def
  by presburger
  then have ⟨angle-c ?Q' C A = ?γ⟩
  by (smt (verit, ccfv-SIG) ⟨A ≠ C ∧ Q ∉ line A C⟩ ⟨angle-c Q A C = angle-c
    B A C / 3⟩
    complex-rotation.ang-eq-theta angle-c-commute angle-symmetry angle-symmetry-eq-imp
    assms(7))

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    axial-symmetry-eq axial-symmetry-eq-line f10 f13 complex-rotation.img-eqI
    neq-0 neq-pi)
  then have ⟨angle-c A C Q = ?γ⟩
    using assms(7) by blast
  then have ⟨angle-c ?Q' C Q = |2*?γ|⟩
    by (metis ⟨A ≠ C ∧ Q ∉ line A C⟩ ⟨angle-c (axial-symmetry A C Q) C A =
    angle-c A C B / 3⟩
    ⟨angle-c Q A C = angle-c B A C / 3⟩ angle-c-neq0 angle-sum-symmetry
    angle-symmetry-eq
    axial-symmetry-eq involution-symmetry mult-2-right mult-commute-abs neg-equal-0-iff-equal
    neq-0 zero-canonical)
  then have f18:⟨rot3.r ?Q' = Q⟩
    using img-r-sym[of C A ?Q]
    by (metis ⟨A ≠ C ∧ Q ∉ line A C⟩ axial-symmetry-dist1
    axial-symmetry-eq canon-ang-cos canon-ang-sin cis.code
    complex-rotation.img-eqI complex-rotation.r-def)
  have Q-inv:⟨rot3.r (rot1.r Q) = Q⟩
    using f17 f18 by presburger
  let ?R' = ⟨axial-symmetry A B R⟩
  have ⟨angle-c B A R = ?α⟩
    by (simp add: assms(3))
  then have ⟨angle-c R A B = -?α⟩
    by (metis angle-c-commute canon-ang-uminus-pi neq-pi pi-canonical)
  have ⟨angle-c R B A = ?β⟩
    by (metis ⟨angle-c (axial-symmetry B C P) B C = angle-c C B A / 3⟩ an-
    gle-c-commute
    angle-c-commute-pi assms(1) assms(2) collinear-sin-neq-0 collinear-sym1
    minus-divide-left sin-pi)
  have ⟨A ≠ B ∧ R ∉ line A B⟩
    by (metis (no-types, opaque-lifting) ⟨angle-c Q A C = angle-c B A C / 3⟩
    ⟨angle-c R A B = -
    (angle-c B A C / 3)⟩ angle-c-commute angle-c-commute-pi angle-c-neq0 assms(2)
    collinear-angle
    collinear-sym2' diff-zero divide-eq-0-iff line-def mem-Collect-eq minus-diff-eq
    neg-equal-zero
    neq-0 zero-canonical)
  then have f17:⟨rot2.r R = ?R'⟩
    using img-r-sym[of B A R]
    using ⟨angle-c R B A = angle-c C B A / 3⟩ axial-symmetry-eq
    cis.code collinear-sym2 line-def complex-rotation.r-def by auto
  then have ⟨angle-c ?R' A B = ?α⟩
    by (metis ⟨A ≠ B ∧ R ∉ line A B⟩ ⟨angle-c R A B = - (angle-c B A C / 3)⟩
    ⟨angle-c R B A = angle-c C B A / 3⟩ add.inverse-inverse add.inverse-neutral
    angle-c-neq0
    angle-symmetry-eq neq-0 zero-canonical)
  then have f18:⟨rot1.r ?R' = R⟩
    using img-r-sym[of A B ?R']
    by (metis ⟨A ≠ B ∧ R ∉ line A B⟩ axial-symmetry-eq canon-ang-cos canon-ang-sin
    cis.code

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    involution-symmetry line-is-inv complex-rotation.r-def z2-inv)
  have R-inv:⟨rot1.r (rot2.r R) = R⟩
    using f17 f18 by presburger
  let ?a1 = ⟨cis(2*?α)⟩ let ?b1 = ⟨A*(1-?a1)⟩ let ?a2 = ⟨cis(2*?β)⟩ let ?b2 = ⟨B*(1-?a2)⟩
  let ?a3 = ⟨cis(2*?γ)⟩ let ?b3 = ⟨C*(1-?a3)⟩
  have possi-abc:⟨⌊?α+?β+?γ⌋ = pi/3 ∨ ⌊?α+?β+?γ⌋ = -pi/3 ∨ ⌊?α+?β+?γ⌋
= pi⟩
  proof -
    have ⟨⌊?α+?β+?γ⌋ ∈ {-pi <.. pi}⟩
      by (simp add: canon-ang(1) canon-ang(2))
    then have ⟨3*⌊?α+?β+?γ⌋ ∈ {-3*pi <.. 3*pi}⟩
      by (auto)
    then have ⟨⌊⌊?α+?β+?γ⌋+⌊?α+?β+?γ⌋+⌊?α+?β+?γ⌋⌋ = ⌊?α+?β+?γ+?α+?β+?γ+?α+?β+?γ⌋⟩
      by (smt (verit, ccfv-SIG) add.commute arg-cis canon-ang-arg
        canon-ang-sum distrib-left is-num-normalize(1) mult-2)
    then have ⟨⌊3*⌊?α+?β+?γ⌋⌋ = pi⟩
      using eq-pi by argo
    then have ⟨3*⌊?α+?β+?γ⌋ = 3*pi ∨ 3*⌊?α+?β+?γ⌋ = -pi ∨ 3*⌊?α+?β+?γ⌋ = pi⟩
      by (smt (verit, del-insts) canon-ang(1) canon-ang(2) canon-ang-id
        canon-ang-minus-3pi-minus-pi canon-ang-pi-3pi)
    then show ?thesis
      by force
  qed
  have jj:⟨B ∉ line C A⟩ ⟨Q ∉ line A C⟩ ⟨R ∉ line A B⟩
    using assms f3' ⟨A ≠ C ∧ Q ∉ line A C⟩ ⟨A ≠ B ∧ R ∉ line A B⟩
    by (simp-all add: collinear-sym1 collinear-sym2 line-def)
  then have inf-pi3:⟨abs ?α < pi/3⟩ ⟨abs ?β < pi/3⟩ ⟨abs ?γ < pi/3⟩
    using ang-inf-pi3[of B C A ?α] ang-inf-pi3[of C A B ?β] ang-inf-pi3[of A B C
?γ]
    by (auto simp add: collinear-sym2 collinear-sym1 assms line-def)
  then have ⟨abs (?α+?β+?γ) < pi⟩
    by argo
  have j1:⟨⌊?α+?β+?γ⌋ = pi/3 ⟹ ?a1*?a2*?a3 = root3⟩
  proof -
    assume hh:⟨⌊?α+?β+?γ⌋ = pi/3⟩
    have f20:⟨cis (2 * (angle-c B A C / 3)) * cis (2 * (angle-c C B A / 3)) * cis
(2 * (angle-c A C B / 3))
= cis (2*(?α+?β+?γ))⟩
      by (simp add: cis-mult)
    then have f21:⟨cis (2*(⌊?α+?β+?γ⌋)) = cis (2*(?α+?β+?γ))⟩
      by (smt (verit, ccfv-threshold) canon-ang-cos canon-ang-sin canon-ang-sum
cis.code)
    with hh show ?thesis
      by (metis (no-types, opaque-lifting) f21 f20 times-divide-eq-right root3-def)
  qed
  have ⟨rot1.r ◦ rot1.r = complex-rotation.r A (4*?α)⟩
    using composed-rotations-same-center by auto
  have g10:⟨((rot1.r ◦ rot1.r ◦ rot1.r) ◦ (rot2.r ◦ rot2.r ◦ rot2.r) ◦ (rot3.r ◦ rot3.r ◦ rot3.r))
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      complex-rotation.r A (6*?α) ∘ complex-rotation.r B (6*?β) ∘ com-
plex-rotation.r C (6*?γ)
    using composed-rotations-same-center by auto
    have f26:⟨|6*?α + 6*?β + 6*?γ| = 0⟩
      by (smt (verit, ccfv-SIG) canon-ang-sum eq-pi two-pi-canonical)
    also have ⟨|6*?γ| ≠ 0⟩
    proof (rule ccontr)
      assume ⟨¬ |6 * (angle-c A C B / 3)| ≠ 0⟩
      then obtain k::int where ⟨6 * |(angle-c A C B / 3)| = 2*k*pi⟩
        by (metis add.commute add.inverse-neutral add-0 canon-ang(3)
            diff-conv-add-uminus f10 f13 of-int-mult of-int-numeral)
      then have ⟨3*|(angle-c A C B / 3)| = k*pi⟩
        by force
      then show False
        by (metis (no-types, opaque-lifting) assms(1) collinear-sin-neq-0 divide-eq-eq-numeral1(1)
            f10 f13 mult-1 sin-kpi times-divide-eq-left zero-neq-numeral)
    qed
    then have f24:⟨|6*?α + 6*?β| ≠ 0⟩
      by (metis add.commute add.right-neutral arg-cis calculation
            canon-ang-cos canon-ang-sin canon-ang-sum cis.code)
    let ?A1:=⟨(A*(1-cis(6*?α)) + B*cis(6*?α)*(1-cis(6*?β)))/(1-cis(6*?β+6*?α))⟩
    have g11:⟨complex-rotation.r A (6*?α) ∘ complex-rotation.r B (6*?β) ∘ com-
plex-rotation.r C (6*?γ) =
      complex-rotation.r ?A1 (6*?α + 6*?β) ∘ complex-rotation.r C (6*?γ) ⟩
      using composed-rotations[of ⟨6*?α⟩ ⟨6*?β⟩ A B, OF f24]
      by argo
    then have f27:⟨complex-rotation.r ?A1 (6*?α + 6*?β) ∘ complex-rotation.r C
(6*?γ) = (λz. z + (C -
      (A * (1 - cis (6 * (angle-c B A C / 3))) + B * cis (6 * (angle-c B A C /
3))) * (1 - cis (6 * (angle-c C B A / 3)))) /
      (1 - cis (6 * (angle-c C B A / 3) + 6 * (angle-c B A C / 3))) *
      (cis (6 * (angle-c B A C / 3) + 6 * (angle-c C B A / 3)) - 1)⟩ (is ⟨?l =
(λz. z + ?A2)⟩)
      using composed-rotation-is-trans[of ⟨6*?α + 6*?β⟩ ⟨6*?γ⟩ ?A1 C, OF f26]
      by auto
    have f30:⟨2*angle-c A C B = 6*?γ⟩
      by auto
    have ⟨axial-symmetry C B A = complex-rotation.r C (|2*angle-c A C B|) A⟩
      using img-r-sym collinear-sym1 f2 jj(1) line-def by auto
    then have first:⟨complex-rotation.r C (6*?γ) A = axial-symmetry C B A⟩ (is
⟨?rr = ?axA⟩)
      using canon-ang-cos canon-ang-sin cis.code f30 complex-rotation.r-def by pres-
burger
    have f31:⟨axial-symmetry B C ?axA = complex-rotation.r B (|2*angle-c ?axA B
C|) ?axA⟩
      by (metis ⟨A ≠ C ∧ Q ∉ line A C⟩ ⟨|6 * (angle-c A C B / 3)| ≠ 0⟩
          ⟨complex-rotation.r C (6 * (angle-c A C B / 3)) A = axial-symmetry C B
A⟩
          complex-rotation.ang-eq-theta angle-c-neq0

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    axial-symmetry-eq f2 img-r-sym involution-symmetry line-is-inv z1-inv)
  have f32:⟨angle-c C B A = 3*?β⟩
  by simp
  then have f33:⟨angle-c ?axA B C = 3*?β⟩
  by (smt (verit) ⟨A ≠ B ∧ R ∉ line A B⟩ ⟨A ≠ C ∧ Q ∉ line A C⟩ ⟨6 *
(angle-c A C B / 3)⟩ ≠ 0⟩
    ⟨complex-rotation.r C (6 * (angle-c A C B / 3)) A = axial-symmetry C B
A⟩
    complex-rotation.ang-eq-theta angle-c-commute angle-c-neq0 angle-symmetry-eq
assms(1)
    axial-symmetry-eq collinear-sin-neq-0 collinear-sym1 f2 line-is-inv sin-pi)
  then have snd:⟨complex-rotation.r B (6*?β) ((axial-symmetry B C A)) = A⟩
  by (smt (verit, ccfv-SIG) ⟨A ≠ C ∧ Q ∉ line A C⟩ ⟨6 * (angle-c A C B / 3)⟩
≠ 0⟩
    ⟨complex-rotation.r C (6 * (angle-c A C B / 3)) A = axial-symmetry C B
A⟩
    complex-rotation.ang-eq-theta angle-c-neq0 axial-symmetry-eq canon-ang-cos
canon-ang-sin
    cis.code f2 img-r-sym involution-symmetry line-is-inv complex-rotation.r-def
z1-inv)
  then have thrd:⟨complex-rotation.r A (6*?α) A = A⟩
  unfolding complex-rotation.r-def by auto
  then have ⟨(complex-rotation.r A (6*?α) ∘ complex-rotation.r B (6*?β) ∘ com-
plex-rotation.r C (6*?γ)) A = A⟩
  apply (simp)
  by (smt (verit, best) axial-symmetry-eq f30 f32 first snd)
  then have g21:⟨((rot1.r ∘ rot1.r ∘ rot1.r) ∘ (rot2.r ∘ rot2.r ∘ rot2.r) ∘ (rot3.r ∘ rot3.r ∘ rot3.r))
= (λz. z + ?A2)⟩
  apply (subst g10) apply (subst g11) apply (subst f27) by (simp add: fun-eq-iff)
  then have ⟨?A2 = 0⟩
  by (metis ⟨(complex-rotation.r A (6 * (angle-c B A C / 3)) ∘
complex-rotation.r B (6 * (angle-c C B A / 3)) ∘
complex-rotation.r C (6 * (angle-c A C B / 3))) A = A⟩
    add-diff-cancel-left' cancel-comm-monoid-add-class.diff-cancel g10)
  then have g22:⟨((rot1.r ∘ rot1.r ∘ rot1.r) ∘ (rot2.r ∘ rot2.r ∘ rot2.r) ∘ (rot3.r ∘ rot3.r ∘ rot3.r))
= id⟩
  using g21 by (auto simp: fun-eq-iff)
  have g20:⟨rot1.r z = ?a1*z + ?b1⟩ ⟨rot2.r z = ?a2*z + ?b2⟩ ⟨rot3.r z = ?a3 *
z + ?b3⟩ for z
  by (auto simp: field-simps complex-rotation.r-def)
  then have ⟨((rot1.r ∘ rot1.r ∘ rot1.r) ∘ (rot2.r ∘ rot2.r ∘ rot2.r) ∘ (rot3.r ∘ rot3.r ∘ rot3.r))
z = (cis (2 * (angle-c B A C / 3)) * cis (2 * (angle-c C B A / 3))
* cis (2 * (angle-c A C B / 3))) ^ 3 * z +
((cis (2 * (angle-c B A C / 3)))^2 + cis (2 * (angle-c B A C / 3)) + 1) * (A
* (1 - cis (2 * (angle-c B A C / 3)))) +
cis (2 * (angle-c B A C / 3)) ^ 3 * ((cis (2 * (angle-c C B A / 3)))^2 + cis (2
* (angle-c C B A / 3)) + 1) *
(B * (1 - cis (2 * (angle-c C B A / 3)))) +
cis (2 * (angle-c B A C / 3)) ^ 3 * cis (2 * (angle-c C B A / 3)) ^ 3 *

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    ((cis (2 * (angle-c A C B / 3)))2 + cis (2 * (angle-c A C B / 3)) + 1) *
    (C * (1 - cis (2 * (angle-c A C B / 3)))) (is (z = ?A3 z) for z
    using equality-for-comp[of rot3.r ?a3 ?b3 rot2.r ?a2 ?b2 rot1.r ?a1 ?b1, OF
g20(3) g20(2) g20(1)]
    by blast
    then have (z = ?A3 z) for z
    by (metis g22 id-apply)
    then have very-imp:((cis (2 * (angle-c B A C / 3)))2 + cis (2 * (angle-c B A
C / 3)) + 1) * (A * (1 - cis (2 * (angle-c B A C / 3)))) +
    cis (2 * (angle-c B A C / 3)) ^ 3 * ((cis (2 * (angle-c C B A / 3)))2 + cis (2
* (angle-c C B A / 3)) + 1) *
    (B * (1 - cis (2 * (angle-c C B A / 3)))) +
    cis (2 * (angle-c B A C / 3)) ^ 3 * cis (2 * (angle-c C B A / 3)) ^ 3 *
    ((cis (2 * (angle-c A C B / 3)))2 + cis (2 * (angle-c A C B / 3)) + 1) *
    (C * (1 - cis (2 * (angle-c A C B / 3)))) = 0)
    by (metis comm-monoid-add-class.add-0 mult-zero-class.mult-zero-right)
    {assume hhh:(?α+?β+?γ) = pi/3}
    have j5:(?a1*?a2 ≠ 1)
    proof(rule ccontr)
      assume (¬ cis (2 * (angle-c B A C / 3)) * cis (2 * (angle-c C B A / 3)) ≠
1)
      then have (?a1*?a2 = cis 0)
      using cis-zero by presburger
      then have (2*(?α+?β)) = 0)
      by (metis arg-cis cis-mult distrib-left zero-canonical)
      then have t:(?α+?β) = pi ∨ (?α+?β) = 0)
      by (smt (verit, del-Insts) canon-ang(1) canon-ang-id canon-ang-sum canon-ang-uminus)
      then have (?α+?β) = 0 ⇒ False)
      by (metis add-0 asms(1) canon-ang-sum collinear-sin-neq-0
          divide-cancel-right f10 f13 hhh sin-pi zero-neq-numeral)
      then have (1/3*(3*?α+3*?β))=pi) using t by(auto simp:algebra-simps)
      then have ((pi-3*?γ)) = (3*?α+3*?β))
      by (smt (verit, ccfv-SIG) canon-ang-diff eq-pi)
      then have ((pi/3-?γ)) = pi)
      by (smt (verit, best) (angle-c B A C / 3 + angle-c C B A / 3) = 0 ⇒
False)
      canon-ang-diff hhh t)
      then have (?γ) ∈ {0<..

```

```

    by(auto simp:field-simps complex-rotation.r-def cis.code)
  have f21:⟨rot2.r (rot3.r P) = ?a2*?a3*P + ?a2*?b3 + ?b2⟩
    apply(subst r-eq-all(2)) apply(subst r-eq-all(3)) by(auto simp:field-simps)

  then have ⟨rot2.r (rot3.r P) = ?a2*?a3*P + ?a2*?b3 + ?b2 ⟷ P*(1-?a2*?a3)
    = ?a2*?b3 + ?b2⟩
    by (smt (verit) add.commute add-diff-cancel-left' add-diff-eq f15
  f16 mult.commute mult.right-neutral right-diff-distrib')
  then have j2:⟨?a2*?a3 ≠ 1 ⟹ P = (?a2*?b3 + ?b2)/(1-?a2*?a3)⟩
    using f21 by(auto simp:field-simps)
  have j4:⟨?a1*?a3 ≠ 1⟩
  proof(rule ccontr)
    assume ¬ cis (2 * (angle-c B A C / 3)) * cis (2 * (angle-c A C B / 3)) ≠
1⟩
    then have ⟨?a1*?a3 = cis 0⟩
      using cis-zero by presburger
    then have ⟨⌊2*(?α+?γ)⌋ = 0⟩
      by (metis arg-cis cis-mult distrib-left zero-canonical)
    then have t:⟨⌊?α+?γ⌋ = pi ∨ ⌊?α+?γ⌋ = 0⟩
    by (smt (verit, del-Insts) canon-ang(1) canon-ang-id canon-ang-sum canon-ang-uminus)
    then have k0:⟨⌊?α+?γ⌋ = 0 ⟹ False⟩
      by (smt (verit, ccfv-SIG) ⟨A ≠ B ∧ R ∉ line A B⟩ ⟨A ≠ C ∧ Q ∉ line A
C⟩
      ⟨⌊angle-c B A C / 3 + angle-c C B A / 3 + angle-c A C B / 3⌋ < pi⟩
      ⟨angle-c (axial-symmetry A B R) A B = angle-c B A C / 3⟩ ⟨angle-c R
A B = - (angle-c B A C / 3)⟩
      ⟨angle-c R B A = angle-c C B A / 3⟩ ang-pos-pos assms(3) canon-ang-id
f2 f30 f32 neq-0 z1-inv)

    then have ⟨⌊1/3*(3*?α+3*?γ)⌋=pi⟩ using t by(auto simp:algebra-simps)
    then have ⟨⌊(pi-⌊3*?β⌋)⌋ = ⌊(3*?α+3*?γ)⌋⟩
      by (smt (verit, ccfv-SIG) canon-ang-diff eq-pi)
    then have ⟨⌊(pi/3-⌊?β⌋)⌋ = pi⟩
      by (smt (verit, best) k0
      canon-ang-diff hhh t)
    then have k1:⟨⌊?β⌋ ∈ {0<..

```

```

    by (metis arg-cis cis-mult distrib-left zero-canonical)
  then have  $t: \langle \lfloor ?\beta + ?\gamma \rfloor = \pi \vee \lfloor ?\beta + ?\gamma \rfloor = 0 \rangle$ 
  by (smt (verit, del-insts) canon-ang(1) canon-ang-id canon-ang-sum canon-ang-uminus)
  then have  $k0: \langle \lfloor ?\beta + ?\gamma \rfloor = 0 \implies \text{False} \rangle$ 
  by (smt (verit, ccfv-SIG)  $\langle A \neq B \wedge R \notin \text{line } A \ B \rangle \langle A \neq C \wedge Q \notin \text{line } A$ 
 $C \rangle$ 
     $\langle \text{angle-c } B \ A \ C \ / \ 3 + \text{angle-c } C \ B \ A \ / \ 3 + \text{angle-c } A \ C \ B \ / \ 3 \rangle < \pi \rangle$ 
     $\langle \text{angle-c } (\text{axial-symmetry } A \ B \ R) \ A \ B = \text{angle-c } B \ A \ C \ / \ 3 \rangle \langle \text{angle-c } R$ 
 $A \ B = -(\text{angle-c } B \ A \ C \ / \ 3) \rangle$ 
     $\langle \text{angle-c } R \ B \ A = \text{angle-c } C \ B \ A \ / \ 3 \rangle \text{ang-pos-pos assms}(3) \text{ canon-ang-id}$ 
 $f2 \ f30 \ f32 \ \text{neq-0 } z1\text{-inv})$ 

  then have  $\langle \lfloor 1/3 * (3 * ?\beta + 3 * ?\gamma) \rfloor = \pi \rangle$  using  $t$  by (auto simp: algebra-simps)
  then have  $\langle \lfloor (\pi - \lfloor 3 * ?\alpha \rfloor) \rfloor = \lfloor (3 * ?\beta + 3 * ?\gamma) \rfloor \rangle$ 
  by (smt (verit, ccfv-SIG) canon-ang-diff eq-pi)
  then have  $\langle \lfloor (\pi/3 - \lfloor ?\alpha \rfloor) \rfloor = \pi \rangle$ 
  by (smt (verit, best)  $k0$ 
    canon-ang-diff hhh  $t$ )
  then have  $k1: \langle \lfloor ?\alpha \rfloor \in \{0 < .. < \pi/3\} \rangle$ 
  by (smt (verit, del-insts) arccos-one-half arcsin-one-half arcsin-plus-arccos
    canon-ang-id divide-nonneg-pos field-sum-of-halves inf-pi3(1) pi-ge-zero)
  then show False
  using  $k1$  add-divide-distrib canon-ang-id  $\langle \lfloor (\pi/3 - \lfloor ?\alpha \rfloor) \rfloor = \pi \rangle$  by auto
qed
have  $f21': \langle \text{rot3.r } (\text{rot1.r } Q) = ?a3 * ?a1 * Q + ?a3 * ?b1 + ?b3 \rangle$ 
  apply (subst r-eq-all(1)) apply (subst r-eq-all(3)) by (auto simp: field-simps)
have  $f21'': \langle \text{rot1.r } (\text{rot2.r } R) = ?a1 * ?a2 * R + ?a1 * ?b2 + ?b1 \rangle$ 
  apply (subst r-eq-all(1)) apply (subst r-eq-all(2)) by (auto simp: field-simps)
have  $jjj: \langle ?a1 \neq 0 \wedge ?a2 \neq 0 \wedge ?a3 \neq 0 \rangle$ 
  using cis-neq-zero by blast
then have  $\text{eq-j}: \langle ?a1 * ?a2 * ?a3 = \text{root3} \rangle$ 
  using  $j1$  hhh by blast
then have  $P\text{-def}: \langle P = (?a2 * ?b3 + ?b2) / (1 - ?a2 * ?a3) \rangle$  (is  $\langle := ?P \rangle$ )
  using  $j2 \ j3$  by blast
have  $n1: \langle 1 - ?a2 * ?a3 = (?a1 - \text{root3}) / ?a1 \rangle$ 
  by (smt (verit, ccfv-threshold) cis-divide cis-times-cis-opposite cis-zero
    diff-divide-distrib eq-j minus-real-def mult-1 times-divide-eq-left)
then have  $P\text{-last}: \langle P = ?a1 * (?a2 * ?b3 + ?b2) / (?a1 - \text{root3}) \rangle$ 
  using  $P\text{-def } jjj$  by (simp add: )
then have  $Q\text{-def}: \langle Q = (?a3 * ?b1 + ?b3) / (1 - ?a3 * ?a1) \rangle$  (is  $\langle := ?Q \rangle$ )
  using  $f21' \ j4 \ Q\text{-inv}$  by (auto simp: field-simps)
have  $n2: \langle 1 - ?a3 * ?a1 = (?a2 - \text{root3}) / ?a2 \rangle$ 
  by (smt (verit, ccfv-threshold) cis-divide cis-times-cis-opposite cis-zero
    diff-divide-distrib eq-j minus-real-def mult-1 times-divide-eq-left)
then have  $Q\text{-last}: \langle Q = ?a2 * (?a3 * ?b1 + ?b3) / (?a2 - \text{root3}) \rangle$ 
  using  $Q\text{-def } jjj$  by simp
then have  $R\text{-def}: \langle R = (?a1 * ?b2 + ?b1) / (1 - ?a1 * ?a2) \rangle$  (is  $\langle := ?R \rangle$ )
  using  $f21'' \ j5 \ R\text{-inv}$  by (auto simp: field-simps)
have  $n3: \langle 1 - ?a1 * ?a2 = (?a3 - \text{root3}) / ?a3 \rangle$ 

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    by (smt (verit, ccfv-threshold) cis-divide cis-times-cis-opposite cis-zero
        diff-divide-distrib eq-j minus-real-def mult-1 times-divide-eq-left)
  then have R-last:  $\langle R = ?a3 * (?a1 * ?b2 + ?b1) / (?a3 - \text{root3}) \rangle$ 
    using R-def jjj by simp
  have  $\langle ?a1 - \text{root3} \neq 0 \wedge ?a2 - \text{root3} \neq 0 \wedge ?a3 - \text{root3} \neq 0 \rangle$ 
    using eq-j j3 j4 j5 by auto
  have rule-s:  $\langle c \neq 0 \implies a/c + b/c = (a+b)/c \rangle$  for  $a\ b\ c::\text{real}$ 
    by argo
  define j' where
    defs:  $\langle j' \equiv \text{root3} \rangle$ 
  have simp-rules-for-eq:  $\langle P = (?a2 * ?b3 + ?b2) / (1 - ?a2 * ?a3) \rangle$   $\langle R = (?a1 * ?b2 + ?b1) / (1 - ?a1 * ?a2) \rangle$ 
     $\langle Q = (?a3 * ?b1 + ?b3) / (1 - ?a1 * ?a3) \rangle$   $\langle 1 - ?a1 * ?a2 \neq 0 \wedge 1 - ?a2 * ?a3 \neq 0 \wedge 1 - ?a1 * ?a3 \neq 0 \rangle$ 
     $\langle ?a1 * ?a2 * ?a3 = j' \rangle$ 
     $\langle 1 + j' + j'^2 = 0 \rangle$   $\langle ((1 - ?a1 * ?a3) * ((1 - ?a1 * ?a2) * (1 - ?a2 * ?a3))) \neq 0 \rangle$ 
     $\langle j'^3 = 1 \rangle$   $\langle j' * j' * j' = 1 \rangle$   $\langle ?a1 \neq 0 \rangle$   $\langle ?a2 \neq 0 \rangle$   $\langle ?a3 \neq 0 \rangle$ 
     $\langle (1 - ?a1 * ?a2) * ?a3 = (?a3 - j') \rangle$   $\langle (1 - ?a2 * ?a3) * ?a1 = (?a1 - j') \rangle$   $\langle (1 - ?a1 * ?a3) * ?a2 = (?a2 - j') \rangle$ 
    using Q-def P-def R-def eq-j j3 jjj j4 j5 root3-eq-0 root-unity-3 n1 n2 n3
    by (auto simp add: mult.commute power3-eq-cube defs root3-def)
  have graal:  $\langle (?a1^2 + ?a1 + 1) * ?b1 + ?a1^3 * (?a2^2 + ?a2 + 1) * ?b2 + ?a1^3 * ?a2^3 * (?a3^2 + ?a3 + 1) * ?b3 \rangle$ 
     $= \langle -j' * ?a1^2 * ?a2 * (?a1 - j') * (?a2 - j') * (?a3 - j') * (?R + j' * ?P + j'^2 * ?Q) \rangle$ 
    using root-unity-carac[of ?a1 ?a2 ?a3 j' ?R ?b2 ?b1 ?P ?b3 ?Q]
    using simp-rules-for-eq defs by (auto simp:)
  then have  $\langle (?a1^2 + ?a1 + 1) * ?b1 + ?a1^3 * (?a2^2 + ?a2 + 1) * ?b2 + ?a1^3 * ?a2^3 * (?a3^2 + ?a3 + 1) * ?b3 \rangle = 0 \rangle$ 
    unfolding defs using very-imp by auto
  then have  $\langle (?R + j' * ?P + j'^2 * ?Q) = 0 \rangle$ 
    using graal simp-rules-for-eq(13) simp-rules-for-eq(14)
    simp-rules-for-eq(15) simp-rules-for-eq(11) simp-rules-for-eq(12) simp-rules-for-eq(4)
  by force
  then have impp:  $\langle R + j' * P + j'^2 * Q = 0 \rangle$ 
    using R-def P-def Q-def
    by (metis simp-rules-for-eq(1) simp-rules-for-eq(2) )
  have neq-all:  $\langle A \neq B \wedge B \neq C \wedge C \neq A \rangle$ 
    using  $\langle A \neq B \wedge R \notin \text{line } A\ B \rangle$   $\langle A \neq C \wedge Q \notin \text{line } A\ C \rangle$  f2 by argo
  have  $\langle R \neq Q \rangle$ 
  proof (rule ccontr)
    assume  $\langle \neg R \neq Q \rangle$ 
    then have q1:  $\langle \text{angle-c } A\ B\ R = \text{angle-c } A\ B\ Q \rangle$ 
       $\langle \text{angle-c } B\ A\ R = \text{angle-c } B\ A\ Q \rangle$   $\langle \text{angle-c } C\ A\ Q = \text{angle-c } C\ A\ R \rangle$ 
       $\langle \text{angle-c } A\ C\ Q = \text{angle-c } A\ C\ R \rangle$ 
      using assms by auto
    then have  $\langle \text{angle-c } B\ A\ R = ?\alpha \rangle$  using assms by auto
    then have  $\langle \text{angle-c } B\ A\ Q = ?\alpha \rangle$ 
      using q1 by auto

```

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then have ⟨angle-c A B Q = angle-c A B R + angle-c R B Q⟩
  using ⟨¬ R ≠ Q⟩ angle-c-neq0 by auto
have ⟨angle-c C A B = - 3 * ?α⟩ using assms
  using ⟨angle-c C A Q = - (angle-c B A C / 3)⟩ by argo
then have ⟨angle-c (axial-symmetry B A R) A B = ?α⟩
  by (metis ⟨angle-c (axial-symmetry A B R) A B = angle-c B A C / 3⟩
axial-symmetry-eq)
have ⟨C - A ≠ 0 ∧ Q - A ≠ 0 ∧ A - R ≠ 0 ∧ A - B ≠ 0⟩
  using ⟨A ≠ C ∧ Q ∉ line A C⟩ ⟨angle-c A B Q = angle-c A B R + angle-c
R B Q⟩
  ⟨angle-c R B A = angle-c C B A / 3⟩ neq-0 q1(1) neq-all by fastforce
then have ⟨angle-c C A B = angle-c C A Q + angle-c Q A R + angle-c R
A B⟩
  using angle-c-sum[of C A Q R] angle-c-sum[of C A R B]
  by (smt (verit) ⟨¬ R ≠ Q⟩ ⟨angle-c C A B = - 3 * (angle-c B A C / 3)⟩
    ⟨angle-c C A Q = - (angle-c B A C / 3)⟩ ⟨angle-c R A B = - (angle-c
B A C / 3)⟩
    angle-c-neq0 canon-ang(1) canon-ang(2) canon-ang-id)
then show False
  using ⟨¬ R ≠ Q⟩
  by (smt (verit, best) ⟨angle-c C A B = - 3 * (angle-c B A C / 3)⟩
    ⟨angle-c C A Q = - (angle-c B A C / 3)⟩ ⟨angle-c R A B = - (angle-c
B A C / 3)⟩
    angle-c-neq0 neq-0 zero-canonical)
qed
then have ⟨cdist R P = cdist R Q ∧ cdist R Q = cdist Q P⟩
  using equilateral-characterization-neg[of R Q P] impp unfolding defs
  by (metis cdist-commute)
then have ?thesis
  using cdist-commute by auto
}note case-pi3=this
then show ?thesis using hhh by auto
qed

theorem Morley-neg:
  assumes ⟨¬ collinear A B C⟩
    ⟨angle-c A B R = angle-c A B C / 3⟩ (is ⟨?abr = ?abc⟩)
    angle-c B A R = angle-c B A C / 3 (is ⟨?bar = ?α⟩)
    angle-c B C P = angle-c B C A / 3 (is ⟨?bcp = ?bca⟩)
    angle-c C B P = angle-c C B A / 3 (is ⟨?cbp = ?β⟩)
    angle-c C A Q = angle-c C A B / 3 (is ⟨?caq = ?cab⟩)
    angle-c A C Q = angle-c A C B / 3 (is ⟨?acq = ?γ⟩)
    and hhh: ⟨angle-c B A C / 3 + angle-c C B A / 3 + angle-c A C B / 3 =
-pi/3⟩
  shows ⟨cdist R P = cdist P Q ∧ cdist Q R = cdist P Q⟩
proof -
  have ⟨angle-c B A C / 3 + angle-c C B A / 3 + angle-c A C B / 3 = pi/3⟩
    using hhh
  by (metis (no-types, opaque-lifting) canon-ang(1) canon-ang-uminus less-eq-real-def

```



```

linorder-not-le minus-add-distrib minus-divide-left mult-le-0-iff
nonzero-mult-div-cancel-left not-numeral-le-zero pi-gt-zero times-divide-eq-right
verit-minus-simplify(4) zero-canonical)
then show ?thesis using Morley-pos[of C B A P Q R]
by (smt (verit, best) Morley-pos angle-c-commute assms(1) assms(2) assms(3)
assms(4) assms(5)
assms(6) assms(7) cdist-commute collinear-sin-neq-0 collinear-sym1 collinear-sym2
sin-pi)
qed

```

theorem Morley:

```

assumes ¬collinear A B C
⟨angle-c A B R = angle-c A B C / 3⟩ (is ⟨?abr = ?abc⟩)
angle-c B A R = angle-c B A C / 3 (is ⟨?bar = ?α⟩)
angle-c B C P = angle-c B C A / 3 (is ⟨?bcp = ?bca⟩)
angle-c C B P = angle-c C B A / 3 (is ⟨?cbp = ?β⟩)
angle-c C A Q = angle-c C A B / 3 (is ⟨?caq = ?cab⟩)
angle-c A C Q = angle-c A C B / 3 (is ⟨?acq = ?γ⟩)
shows ⟨cdist R P = cdist P Q ∧ cdist Q R = cdist P Q⟩

```

proof –

```

{fix A B C γ
assume ABC-nline:⟨A ∉ line B C⟩
and eq-3c:⟨angle-c A C B = 3*γ⟩
then have neq-PI:⟨abs γ < pi/3⟩
proof –
have ⟨angle-c A C B ≠ pi⟩ using ABC-nline(1) unfolding line-def
by (metis angle-c-commute-pi collinear-iff mem-Collect-eq non-collinear-independant)
then have ⟨angle-c A C B ∈ {−pi<..using ang-c-in less-eq-real-def by auto
then show ⟨abs γ < pi/3⟩
using eq-3c by force
qed}note ang-inf-pi3=this
have ⟨|angle-c B A C + angle-c C B A + angle-c A C B| = pi⟩
by (metis Elementary-Complex-Geometry.collinear-def add.commute
angle-sum-triangle-c assms(1) collinear-sym1 collinear-sym2)
have eq-pi:⟨|3*?α + 3*?β + 3*?γ| = pi⟩
using ⟨|angle-c B A C + angle-c C B A + angle-c A C B| = pi⟩ by force
have ⟨A ∉ line B C ∧ B ∉ line A C ∧ C ∉ line A B⟩
using assms(1) unfolding line-def using collinear-sym1 collinear-sym2' by (auto)

```

```

then have f2:⟨abs ?α < pi/3 ∧ abs ?β < pi/3 ∧ abs ?γ < pi/3⟩
using ang-inf-pi3
by (smt (verit, ccfv-SIG) ang-vec-bounded angle-c-def assms(1) collinear-sin-neq-0

```

```

collinear-sym1 collinear-sym2 divide-less-cancel minus-divide-left sin-pi)
have possi-abc:⟨|?α + ?β + ?γ| = pi/3 ∨ |?α + ?β + ?γ| = −pi/3⟩
proof –
have ⟨?α + ?β + ?γ ≤ pi⟩

```

```

    using f2 by (auto simp: field-simps)
  then have  $\langle \lfloor ?\alpha + ?\beta + ?\gamma \rfloor = ?\alpha + ?\beta + ?\gamma \rangle$ 
    using canon-ang-id f2 by fastforce
  then have  $\langle \lfloor 3 * \lfloor ?\alpha + ?\beta + ?\gamma \rfloor \rfloor = \pi \rangle$ 
    using eq-pi by force
  then show ?thesis
    by (smt (verit) add-divide-distrib arccos-minus-one-half arccos-one-half arc-
sin-minus-one-half
arcsin-one-half arcsin-plus-arccos canon-ang-id canon-ang-minus-3pi-minus-pi canon-ang-pi-3pi

eq-pi f2 field-sum-of-halves)
qed
then show ?thesis
  using Morley-neg Morley-pos assms(1) assms(2) assms(3) assms(4) assms(5)
  assms(6) assms(7)
  by meson
qed
end

```

References

- [1] Manuel Eberl *Basic Geometric Properties of Triangles*, Archive of Formal Proofs, December 2015 <https://isa-afp.org/entries/Triangle.html>
- [2] Morley theorem <http://denisevellachemla.eu/Alain-Connes-Theoreme-Morley.pdf>.
- [3] Wiedijk's catalogue "Formalizing 100 Theorems" <https://www.cs.ru.nl/~freek/100/Itappearsatposition121...>