

Lovasz Local Lemma

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Abstract

This entry aims to formalise several useful general techniques for using the *probabilistic method* for combinatorial structures (or discrete spaces more generally). In particular, it focuses on bounding tools, such as the union and complete independence bounds, and the first formalisation of the pivotal Lovász local lemma. The formalisation focuses on the general lemma, however also proves several useful variations, including the more well known symmetric version. Both the original formalisation and several of the variations used dependency graphs, which were formalised using Noschinski’s general directed graph library [2]. Additionally, the entry provides several useful existence lemmas, required at the end of most probabilistic proofs on combinatorial structures. Finally, the entry includes several significant extensions to the existing probability libraries, particularly for conditional probability (such as Bayes theorem) and independent events. The formalisation is primarily based on Alon and Spencer’s textbook [1], as well as Zhao’s course notes [3].

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1 Extensional function extras

Counting lemmas (i.e. reasoning on cardinality) of sets on the extensional function relation

```
theory PiE-Rel-Extras imports Card-Partitions.Card-Partitions
begin
```

1.1 Relations and Extensional Function sets

A number of lemmas to convert between relations and functions for counting purposes. Note, ultimately not needed in this formalisation, but may be of use in the future

```
lemma Range-unfold: Range r = {y.  $\exists x. (x, y) \in r$ }
  <proof>
```

```
definition fun-to-rel:: 'a set  $\Rightarrow$  'b set  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  ('a  $\times$  'b) set where
fun-to-rel A B f  $\equiv$  {(a, b) | a b . a  $\in$  A  $\wedge$  b  $\in$  B  $\wedge$  f a = b}
```

```
definition rel-to-fun:: ('a  $\times$  'b) set  $\Rightarrow$  ('a  $\Rightarrow$  'b) where
rel-to-fun R  $\equiv$   $\lambda a .$  (if a  $\in$  Domain R then (THE b . (a, b)  $\in$  R) else undefined)
```

```
lemma fun-to-relI: a  $\in$  A  $\Longrightarrow$  b  $\in$  B  $\Longrightarrow$  f a = b  $\Longrightarrow$  (a, b)  $\in$  fun-to-rel A B f
  <proof>
```

```
lemma fun-to-rel-alt: fun-to-rel A B f  $\equiv$  {(a, f a) | a b . a  $\in$  A  $\wedge$  f a  $\in$  B}
  <proof>
```

```
lemma fun-to-relI2: a  $\in$  A  $\Longrightarrow$  f a  $\in$  B  $\Longrightarrow$  (a, f a)  $\in$  fun-to-rel A B f
  <proof>
```

lemma *rel-to-fun-in[simp]*: $a \in \text{Domain } R \implies (\text{rel-to-fun } R) a = (\text{THE } b . (a, b) \in R)$

<proof>

lemma *rel-to-fun-undefined[simp]*: $a \notin \text{Domain } R \implies (\text{rel-to-fun } R) a = \text{undefined}$

<proof>

lemma *single-valued-unique-Dom-iff*: $\text{single-valued } R \iff (\forall x \in \text{Domain } R. \exists! y . (x, y) \in R)$

<proof>

lemma *rel-to-fun-range*:

assumes *single-valued* R

assumes $a \in \text{Domain } R$

shows $(\text{THE } b . (a, b) \in R) \in \text{Range } R$

<proof>

lemma *rel-to-fun-extensional*: $\text{single-valued } R \implies \text{rel-to-fun } R \in (\text{Domain } R \rightarrow_E \text{Range } R)$

<proof>

lemma *single-value-fun-to-rel*: $\text{single-valued } (\text{fun-to-rel } A B f)$

<proof>

lemma *fun-to-rel-domain*:

assumes $f \in A \rightarrow_E B$

shows $\text{Domain } (\text{fun-to-rel } A B f) = A$

<proof>

lemma *fun-to-rel-range*:

assumes $f \in A \rightarrow_E B$

shows $\text{Range } (\text{fun-to-rel } A B f) \subseteq B$

<proof>

lemma *rel-to-fun-to-rel*:

assumes $f \in A \rightarrow_E B$

shows $\text{rel-to-fun } (\text{fun-to-rel } A B f) = f$

<proof>

lemma *fun-to-rel-to-fun*:

assumes *single-valued* R

shows $\text{fun-to-rel } (\text{Domain } R) (\text{Range } R) (\text{rel-to-fun } R) = R$

<proof>

lemma *bij-betw-fun-to-rel*:

assumes $f \in A \rightarrow_E B$

shows *bij-betw* $(\lambda a . (a, f a)) A (\text{fun-to-rel } A B f)$

<proof>

lemma *fun-to-rel-indiv-card*:
assumes $f \in A \rightarrow_E B$
shows $\text{card } (\text{fun-to-rel } A \ B \ f) = \text{card } A$
 $\langle \text{proof} \rangle$

lemma *fun-to-rel-inj*:
assumes $C \subseteq A \rightarrow_E B$
shows *inj-on* $(\text{fun-to-rel } A \ B)$ C
 $\langle \text{proof} \rangle$

lemma *fun-to-rel-ss*: $\text{fun-to-rel } A \ B \ f \subseteq A \times B$
 $\langle \text{proof} \rangle$

lemma *card-fun-to-rel*: $C \subseteq A \rightarrow_E B \implies \text{card } C = \text{card } ((\lambda f . \text{fun-to-rel } A \ B \ f) \ ` \ C)$
 $\langle \text{proof} \rangle$

1.2 Cardinality Lemmas

Lemmas to count variations of filtered sets over the extensional function set relation

lemma *card-PiE-filter-range-set*:
assumes $\bigwedge a. a \in A' \implies X \ a \in C$
assumes $A' \subseteq A$
assumes *finite* A
shows $\text{card } \{f \in A \rightarrow_E C . \forall a \in A' . f \ a = X \ a\} = (\text{card } C) \frown (\text{card } A - \text{card } A')$
 $\langle \text{proof} \rangle$

lemma *card-PiE-filter-range-indiv*: $X \ a' \in C \implies a' \in A \implies \text{finite } A \implies \text{card } \{f \in A \rightarrow_E C . f \ a' = X \ a'\} = (\text{card } C) \frown (\text{card } A - 1)$
 $\langle \text{proof} \rangle$

lemma *card-PiE-filter-range-set-const*: $c \in C \implies A' \subseteq A \implies \text{finite } A \implies \text{card } \{f \in A \rightarrow_E C . \forall a \in A' . f \ a = c\} = (\text{card } C) \frown (\text{card } A - \text{card } A')$
 $\langle \text{proof} \rangle$

lemma *card-PiE-filter-range-set-nat*: $c \in \{0..<n\} \implies A' \subseteq A \implies \text{finite } A \implies \text{card } \{f \in A \rightarrow_E \{0..<n\} . \forall a \in A' . f \ a = c\} = n \frown (\text{card } A - \text{card } A')$
 $\langle \text{proof} \rangle$

end

2 Digraph extensions

Extensions to the existing library for directed graphs, basically neighborhood

theory *Digraph-Extensions*
imports

```

    Graph-Theory.Digraph
    Graph-Theory.Pair-Digraph
begin

definition (in pre-digraph) neighborhood :: 'a ⇒ 'a set where
neighborhood u ≡ {v ∈ verts G . dominates G u v}

lemma (in wf-digraph) neighborhood-wf: neighborhood v ⊆ verts G
⟨proof⟩

lemma (in pair-pre-digraph) neighborhood-alt:
neighborhood u = {v ∈ pverts G . (u, v) ∈ parcs G}
⟨proof⟩

lemma (in fin-digraph) neighborhood-finite: finite (neighborhood v)
⟨proof⟩

lemma (in wf-digraph) neighborhood-edge-iff: y ∈ neighborhood x ⟷ (x, y) ∈
arcs-ends G
⟨proof⟩

lemma (in loopfree-digraph) neighborhood-self-not: v ∉ (neighborhood v)
⟨proof⟩

lemma (in nomulti-digraph) inj-on-head-out-arcs: inj-on (head G) (out-arcs G u)
⟨proof⟩

lemma (in nomulti-digraph) out-degree-neighborhood: out-degree G u = card (neighborhood
u)
⟨proof⟩

lemma (in digraph) neighborhood-empty-iff: out-degree G u = 0 ⟷ neighborhood
u = {}
⟨proof⟩

end

```

3 General Event Lemmas

General lemmas for reasoning on events in probability spaces after different operations

```

theory Prob-Events-Extras
  imports
    HOL-Probability.Probability
    PiE-Rel-Extras
begin

context prob-space

```

begin

lemma *prob-sum-Union*:

assumes *measurable*: *finite A* $A \subseteq \text{events}$ *disjoint A*

shows $\text{prob} (\bigcup A) = (\sum_{e \in A} \text{prob} (e))$

<proof>

lemma *events-inter*:

assumes *finite S*

assumes $S \neq \{\}$

shows $(\bigwedge A. A \in S \implies A \in \text{events}) \implies \bigcap S \in \text{events}$

<proof>

lemma *events-union*:

assumes *finite S*

shows $(\bigwedge A. A \in S \implies A \in \text{events}) \implies \bigcup S \in \text{events}$

<proof>

lemma *prob-inter-set-lt-elem*: $A \in \text{events} \implies \text{prob} (A \cap (\bigcap AS)) \leq \text{prob} A$

<proof>

lemma *Inter-event-ss*: $\text{finite } A \implies A \subseteq \text{events} \implies A \neq \{\} \implies \bigcap A \in \text{events}$

<proof>

lemma *prob-inter-ss-lt*:

assumes *finite A*

assumes $A \subseteq \text{events}$

assumes $B \neq \{\}$

assumes $B \subseteq A$

shows $\text{prob} (\bigcap A) \leq \text{prob} (\bigcap B)$

<proof>

lemma *prob-inter-ss-lt-index*:

assumes *finite A*

assumes $F \text{ ' } A \subseteq \text{events}$

assumes $B \neq \{\}$

assumes $B \subseteq A$

shows $\text{prob} (\bigcap (F \text{ ' } A)) \leq \text{prob} (\bigcap (F \text{ ' } B))$

<proof>

lemma *space-compl-double*:

assumes $S \subseteq \text{events}$

shows $((-) (\text{space } M)) \text{ ' } (((-) (\text{space } M)) \text{ ' } S) = S$

<proof>

lemma *bij-betw-compl-sets*:

assumes $S \subseteq \text{events}$

assumes $S' = ((-) (\text{space } M)) \text{ ' } S$

shows *bij-betw* $((-) (\text{space } M)) S' S$

<proof>

lemma *bij-betw-compl-sets-rev*:

assumes $S \subseteq \text{events}$

assumes $S' = ((-) (\text{space } M)) ' S$

shows *bij-betw* $((-) (\text{space } M)) S S'$

<proof>

lemma *prob0-basic-inter*: $A \in \text{events} \implies B \in \text{events} \implies \text{prob } A = 0 \implies \text{prob } (A \cap B) = 0$

<proof>

lemma *prob0-basic-Inter*: $A \in \text{events} \implies B \subseteq \text{events} \implies \text{prob } A = 0 \implies \text{prob } (A \cap (\bigcap B)) = 0$

<proof>

lemma *prob1-basic-inter*: $A \in \text{events} \implies B \in \text{events} \implies \text{prob } A = 1 \implies \text{prob } (A \cap B) = \text{prob } B$

<proof>

lemma *prob1-basic-Inter*:

assumes $A \in \text{events}$ $B \subseteq \text{events}$

assumes $\text{prob } A = 1$

assumes $B \neq \{\}$

assumes *finite* B

shows $\text{prob } (A \cap (\bigcap B)) = \text{prob } (\bigcap B)$

<proof>

lemma *compl-identity*: $A \in \text{events} \implies \text{space } M - (\text{space } M - A) = A$

<proof>

lemma *prob-addition-rule*: $A \in \text{events} \implies B \in \text{events} \implies$

$\text{prob } (A \cup B) = \text{prob } A + \text{prob } B - \text{prob } (A \cap B)$

<proof>

lemma *compl-subset-in-events*: $S \subseteq \text{events} \implies (-) (\text{space } M) ' S \subseteq \text{events}$

<proof>

lemma *prob-compl-diff-inter*: $A \in \text{events} \implies B \in \text{events} \implies$

$\text{prob } (A \cap (\text{space } M - B)) = \text{prob } A - \text{prob } (A \cap B)$

<proof>

lemma *bij-betw-prod-prob*: *bij-betw* $f A B \implies (\prod_{b \in B}. \text{prob } b) = (\prod_{a \in A}. \text{prob } (f a))$

<proof>

definition *event-compl* :: 'a set \Rightarrow 'a set **where**

event-compl $A \equiv \text{space } M - A$

lemma compl-Union: $A \neq \{\}$ \implies $\text{space } M - (\bigcup A) = (\bigcap a \in A . (\text{space } M - a))$
 ⟨proof⟩

lemma compl-Union-fn: $A \neq \{\}$ \implies $\text{space } M - (\bigcup (F \cdot A)) = (\bigcap a \in A . (\text{space } M - F a))$
 ⟨proof⟩

end

Reasoning on the probability of function sets

lemma card-PiE-val-ss-eq:

assumes *finite* A
assumes $b \in B$
assumes $d \subseteq A$
assumes $B \neq \{\}$
assumes *finite* B
shows $\text{card } \{f \in (A \rightarrow_E B) . (\forall v \in d . f v = b)\} / \text{card } (A \rightarrow_E B) = 1 / ((\text{card } B) \text{ powi } (\text{card } d))$
 (is $\text{card } \{f \in ?C . (\forall v \in d . f v = b)\} / \text{card } ?C = 1 / ((\text{card } B) \text{ powi } (\text{card } d))$)
 ⟨proof⟩

lemma card-PiE-val-indiv-eq:

assumes *finite* A
assumes $b \in B$
assumes $d \in A$
assumes $B \neq \{\}$
assumes *finite* B
shows $\text{card } \{f \in (A \rightarrow_E B) . f d = b\} / \text{card } (A \rightarrow_E B) = 1 / (\text{card } B)$
 (is $\text{card } \{f \in ?C . f d = b\} / \text{card } ?C = 1 / (\text{card } B)$)
 ⟨proof⟩

lemma prob-uniform-ex-fun-space:

assumes *finite* A
assumes $b \in B$
assumes $d \subseteq A$
assumes $B \neq \{\}$
assumes $A \neq \{\}$
assumes *finite* B
shows $\text{prob-space.prob } (\text{uniform-count-measure } (A \rightarrow_E B)) \{f \in (A \rightarrow_E B) . (\forall v \in d . f v = b)\} =$
 $1 / ((\text{card } B) \text{ powi } (\text{card } d))$
 ⟨proof⟩

proposition integrable-uniform-count-measure-finite:

fixes $g :: 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$
shows *finite* $A \implies$ *integrable* $(\text{uniform-count-measure } A) g$
 ⟨proof⟩

end

4 Conditional Probability Library Extensions

```
theory Cond-Prob-Extensions
  imports
    Prob-Events-Extras
    Design-Theory.Multisets-Extras
begin
```

4.1 Miscellaneous Set and List Lemmas

```
lemma nth-image-tl:
  assumes  $xs \neq []$ 
  shows  $nth\ xs\ ' \{1..<length\ xs\} = set(tl\ xs)$ 
<proof>
```

```
lemma exists-list-card:
  assumes finite S
  obtains xs where  $set\ xs = S$  and  $length\ xs = card\ S$ 
<proof>
```

```
lemma bij-betw-inter-empty:
  assumes bij-betw f A B
  assumes  $A' \subseteq A$ 
  assumes  $A'' \subseteq A$ 
  assumes  $A' \cap A'' = \{\}$ 
  shows  $f\ ' \ A' \cap f\ ' \ A'' = \{\}$ 
<proof>
```

```
lemma bij-betw-image-comp-eq:
  assumes bij-betw g T S
  shows  $(F \circ g)\ ' \ T = F\ ' \ S$ 
<proof>
```

```
lemma prod-card-image-set-eq:
  assumes bij-betw f  $\{0..<card\ S\}$  S
  assumes finite S
  shows  $(\prod i \in \{n..<(card\ S)\} . g\ (f\ i)) = (\prod i \in f\ ' \ \{n..<card\ S\} . g\ i)$ 
<proof>
```

```
lemma set-take-distinct-elem-not:
  assumes distinct xs
  assumes  $i < length\ xs$ 
  shows  $xs\ !\ i \notin set\ (take\ i\ xs)$ 
<proof>
```

4.2 Conditional Probability Basics

```
context prob-space
begin
```

Abbreviation to mirror mathematical notations

abbreviation *cond-prob-ev* :: 'a set \Rightarrow 'a set \Rightarrow real ($\mathcal{P}'(- | -')$) **where**
 $\mathcal{P}(B | A) \equiv \mathcal{P}(x \text{ in } M. (x \in B) | (x \in A))$

lemma *cond-prob-inter*: $\mathcal{P}(B | A) = \mathcal{P}(\omega \text{ in } M. (\omega \in B \cap A)) / \mathcal{P}(\omega \text{ in } M. (\omega \in A))$
<proof>

lemma *cond-prob-ev-def*:
assumes $A \in \text{events } B \in \text{events}$
shows $\mathcal{P}(B | A) = \text{prob } (A \cap B) / \text{prob } A$
<proof>

lemma *measurable-in-ev*:
assumes $A \in \text{events}$
shows $\text{Measurable.pred } M (\lambda x . x \in A)$
<proof>

lemma *measure-uniform-measure-eq-cond-prob-ev*:
assumes $A \in \text{events } B \in \text{events}$
shows $\mathcal{P}(A | B) = \mathcal{P}(x \text{ in uniform-measure } M \{x \in \text{space } M. x \in B\}. x \in A)$
<proof>

lemma *measure-uniform-measure-eq-cond-prob-ev2*:
assumes $A \in \text{events } B \in \text{events}$
shows $\mathcal{P}(A | B) = \text{measure } (\text{uniform-measure } M \{x \in \text{space } M. x \in B\}) A$
<proof>

lemma *measure-uniform-measure-eq-cond-prob-ev3*:
assumes $A \in \text{events } B \in \text{events}$
shows $\mathcal{P}(A | B) = \text{measure } (\text{uniform-measure } M B) A$
<proof>

lemma *prob-space-cond-prob-uniform*:
assumes $\text{prob } (\{x \in \text{space } M. Q x\}) > 0$
shows $\text{prob-space } (\text{uniform-measure } M \{x \in \text{space } M. Q x\})$
<proof>

lemma *prob-space-cond-prob-event*:
assumes $\text{prob } B > 0$
shows $\text{prob-space } (\text{uniform-measure } M B)$
<proof>

Note this case shouldn't be used. Conditional probability should have > 0 assumption

lemma *cond-prob-empty*: $\mathcal{P}(B | \{\}) = 0$
<proof>

lemma *cond-prob-space*: $\mathcal{P}(A | \text{space } M) = \mathcal{P}(w \text{ in } M . w \in A)$

<proof>

lemma *cond-prob-space-ev*: **assumes** $A \in \text{events}$ **shows** $\mathcal{P}(A \mid \text{space } M) = \text{prob } A$

<proof>

lemma *cond-prob-UNIV*: $\mathcal{P}(A \mid \text{UNIV}) = \mathcal{P}(w \text{ in } M . w \in A)$

<proof>

lemma *cond-prob-UNIV-ev*: $A \in \text{events} \implies \mathcal{P}(A \mid \text{UNIV}) = \text{prob } A$

<proof>

lemma *cond-prob-neg*:

assumes $A \in \text{events}$ $B \in \text{events}$

assumes $\text{prob } A > 0$

shows $\mathcal{P}(\text{space } M - B \mid A) = 1 - \mathcal{P}(B \mid A)$

<proof>

4.3 Bayes Theorem

lemma *prob-intersect-A*:

assumes $A \in \text{events}$ $B \in \text{events}$

shows $\text{prob } (A \cap B) = \text{prob } A * \mathcal{P}(B \mid A)$

<proof>

lemma *prob-intersect-B*:

assumes $A \in \text{events}$ $B \in \text{events}$

shows $\text{prob } (A \cap B) = \text{prob } B * \mathcal{P}(A \mid B)$

<proof>

theorem *Bayes-theorem*:

assumes $A \in \text{events}$ $B \in \text{events}$

shows $\text{prob } B * \mathcal{P}(A \mid B) = \text{prob } A * \mathcal{P}(B \mid A)$

<proof>

corollary *Bayes-theorem-div*:

assumes $A \in \text{events}$ $B \in \text{events}$

shows $\mathcal{P}(A \mid B) = (\text{prob } A * \mathcal{P}(B \mid A)) / (\text{prob } B)$

<proof>

lemma *cond-prob-dual-intersect*:

assumes $A \in \text{events}$ $B \in \text{events}$ $C \in \text{events}$

assumes $\text{prob } C \neq 0$

shows $\mathcal{P}(A \mid (B \cap C)) = \mathcal{P}(A \cap B \mid C) / \mathcal{P}(B \mid C)$ (is ?LHS = ?RHS)

<proof>

lemma *cond-prob-ev-double*:

assumes $A \in \text{events}$ $B \in \text{events}$ $C \in \text{events}$

assumes $\text{prob } C > 0$
shows $\mathcal{P}(x \text{ in } (\text{uniform-measure } M \ C). (x \in A) \mid (x \in B)) = \mathcal{P}(A \mid (B \cap C))$
 <proof>

lemma *cond-prob-inter-set-lt*:
assumes $A \in \text{events } B \in \text{events } AS \subseteq \text{events}$
assumes *finite AS*
shows $\mathcal{P}((A \cap (\bigcap AS)) \mid B) \leq \mathcal{P}(A \mid B)$ (is ?LHS \leq ?RHS)
 <proof>

4.4 Conditional Probability Multiplication Rule

Many list and indexed variations of this lemma

lemma *prob-cond-Inter-List*:
assumes $xs \neq []$
assumes $\bigwedge A. A \in \text{set } xs \implies A \in \text{events}$
shows $\text{prob } (\bigcap (\text{set } xs)) = \text{prob } (\text{hd } xs) * (\prod i = 1..<(\text{length } xs) . \mathcal{P}((xs ! i) \mid (\bigcap (\text{set } (\text{take } i \ xs))))))$
 <proof>

lemma *prob-cond-Inter-index*:
fixes $n :: \text{nat}$
assumes $n > 0$
assumes $F \text{ ' } \{0..<n\} \subseteq \text{events}$
shows $\text{prob } (\bigcap (F \text{ ' } \{0..<n\})) = \text{prob } (F \ 0) * (\prod i \in \{1..<n\} . \mathcal{P}(F \ i \mid (\bigcap (F \text{ ' } \{0..<i\}))))$
 <proof>

lemma *prob-cond-Inter-index-compl*:
fixes $n :: \text{nat}$
assumes $n > 0$
assumes $F \text{ ' } \{0..<n\} \subseteq \text{events}$
shows $\text{prob } (\bigcap x \in \{0..<n\} . \text{space } M - F \ x) = \text{prob } (\text{space } M - F \ 0) * (\prod i \in \{1..<n\} . \mathcal{P}(\text{space } M - F \ i \mid (\bigcap j \in \{0..<i\} . \text{space } M - F \ j)))$
 <proof>

lemma *prob-cond-Inter-take-cond*:
assumes $xs \neq []$
assumes $\text{set } xs \subseteq \text{events}$
assumes $S \subseteq \text{events}$
assumes $S \neq \{\}$
assumes *finite S*
assumes $\text{prob } (\bigcap S) > 0$
shows $\mathcal{P}((\bigcap (\text{set } xs)) \mid (\bigcap S)) = (\prod i = 0..<(\text{length } xs) . \mathcal{P}((xs ! i) \mid (\bigcap (\text{set } (\text{take } i \ xs) \cup S))))$
 <proof>

lemma *prob-cond-Inter-index-cond-set:*

fixes $n :: nat$
assumes $n > 0$
assumes *finite* E
assumes $E \neq \{\}$
assumes $E \subseteq events$
assumes $F \text{ ' } \{0..<n\} \subseteq events$
assumes $prob (\bigcap E) > 0$
shows $\mathcal{P}((\bigcap (F \text{ ' } \{0..<n\})) \mid (\bigcap E)) = (\prod i \in \{0..<n\}. \mathcal{P}(F i \mid (\bigcap ((F \text{ ' } \{0..<i\}) \cup E))))$
<proof>

lemma *prob-cond-Inter-index-cond-compl-set:*

fixes $n :: nat$
assumes $n > 0$
assumes *finite* E
assumes $E \neq \{\}$
assumes $E \subseteq events$
assumes $F \text{ ' } \{0..<n\} \subseteq events$
assumes $prob (\bigcap E) > 0$
shows $\mathcal{P}((\bigcap ((-) (space M) \text{ ' } F \text{ ' } \{0..<n\})) \mid (\bigcap E)) = (\prod i = 0..<n. \mathcal{P}((space M - F i) \mid (\bigcap ((-) (space M) \text{ ' } F \text{ ' } \{0..<i\} \cup E))))$
<proof>

lemma *prob-cond-Inter-index-cond:*

fixes $n :: nat$
assumes $n > 0$
assumes $n < m$
assumes $F \text{ ' } \{0..<m\} \subseteq events$
assumes $prob (\bigcap j \in \{n..<m\}. F j) > 0$
shows $\mathcal{P}((\bigcap (F \text{ ' } \{0..<n\})) \mid (\bigcap j \in \{n..<m\}. F j)) = (\prod i \in \{0..<n\}. \mathcal{P}(F i \mid (\bigcap ((F \text{ ' } \{0..<i\}) \cup (F \text{ ' } \{n..<m\}))))$
<proof>

lemma *prob-cond-Inter-index-cond-compl:*

fixes $n :: nat$
assumes $n > 0$
assumes $n < m$
assumes $F \text{ ' } \{0..<m\} \subseteq events$
assumes $prob (\bigcap j \in \{n..<m\}. F j) > 0$
shows $\mathcal{P}((\bigcap ((-) (space M) \text{ ' } F \text{ ' } \{0..<n\})) \mid (\bigcap (F \text{ ' } \{n..<m\}))) = (\prod i = 0..<n. \mathcal{P}((space M - F i) \mid (\bigcap ((-) (space M) \text{ ' } F \text{ ' } \{0..<i\} \cup (F \text{ ' } \{n..<m\}))))$
<proof>

lemma *prob-cond-Inter-take-cond-neg:*

assumes $xs \neq []$
assumes *set* $xs \subseteq events$

assumes $S \subseteq \text{events}$
assumes $S \neq \{\}$
assumes *finite* S
assumes $\text{prob} (\bigcap S) > 0$
shows $\mathcal{P}((\bigcap((-) (\text{space } M) \text{ ' } (\text{set } xs))) \mid (\bigcap S)) =$
 $(\prod_{i = 0..<(\text{length } xs)} . \mathcal{P}((\text{space } M - xs ! i) \mid (\bigcap((-) (\text{space } M) \text{ ' } (\text{set } (\text{take}$
 $i \text{ } xs \text{ })) \cup S))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-Inter-List-Index:*

assumes $xs \neq []$
assumes $\text{set } xs \subseteq \text{events}$
shows $\text{prob} (\bigcap (\text{set } xs)) = \text{prob} (\text{hd } xs) * (\prod_{i = 1..<(\text{length } xs)} .$
 $\mathcal{P}((xs ! i) \mid (\bigcap_{j \in \{0..<i\}} . xs ! j)))$
 $\langle \text{proof} \rangle$

lemma *obtains-prob-cond-Inter-index:*

assumes $S \neq \{\}$
assumes $S \subseteq \text{events}$
assumes *finite* S
obtains xs **where** $\text{set } xs = S$ **and** $\text{length } xs = \text{card } S$ **and**
 $\text{prob} (\bigcap S) = \text{prob} (\text{hd } xs) * (\prod_{i = 1..<(\text{length } xs)} . \mathcal{P}((xs ! i) \mid (\bigcap_{j \in \{0..<i\}}$
 $. xs ! j)))$
 $\langle \text{proof} \rangle$

lemma *obtain-list-index:*

assumes *bij-betw* $g \{0..<\text{card } S\} S$
assumes *finite* S
obtains xs **where** $\text{set } xs = S$ **and** $\bigwedge i . i \in \{0..<\text{card } S\} \implies g \ i = xs ! i$ **and**
distinct xs
 $\langle \text{proof} \rangle$

lemma *prob-cond-inter-fn:*

assumes *bij-betw* $g \{0..<\text{card } S\} S$
assumes *finite* S
assumes $S \neq \{\}$
assumes $S \subseteq \text{events}$
shows $\text{prob} (\bigcap S) = \text{prob} (g \ 0) * (\prod_{i \in \{1..<(\text{card } S)\}} . \mathcal{P}(g \ i \mid (\bigcap (g \text{ ' } \{0..<i\}))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-inter-obtain-fn:*

assumes $S \neq \{\}$
assumes $S \subseteq \text{events}$
assumes *finite* S
obtains f **where** *bij-betw* $f \{0..<\text{card } S\} S$ **and**
 $\text{prob} (\bigcap S) = \text{prob} (f \ 0) * (\prod_{i \in \{1..<(\text{card } S)\}} . \mathcal{P}(f \ i \mid (\bigcap (f \text{ ' } \{0..<i\}))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-inter-obtain-fn-compl:*

assumes $S \neq \{\}$
assumes $S \subseteq \text{events}$
assumes *finite* S
obtains f **where** *bij-betw* $f \{0..<\text{card } S\} S$ **and** $\text{prob} (\bigcap ((-) (\text{space } M) ' S))$
 $=$
 $\text{prob} (\text{space } M - f 0) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(\text{space } M - f i \mid (\bigcap ((-) (\text{space } M) ' f ' \{0..<i\}))))$
<proof>

lemma *prob-cond-Inter-index-cond-fn:*

assumes $I \neq \{\}$
assumes *finite* I
assumes *finite* E
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F ' I \subseteq \text{events}$
assumes $\text{prob} (\bigcap E) > 0$
assumes *bb: bij-betw* $g \{0..<\text{card } I\} I$
shows $\mathcal{P}((\bigcap (F ' g ' \{0..<\text{card } I\})) \mid (\bigcap E)) =$
 $(\prod i \in \{0..<\text{card } I\} . \mathcal{P}(F (g i) \mid (\bigcap ((F ' g ' \{0..<i\}) \cup E))))$
<proof>

lemma *prob-cond-Inter-index-cond-obtains:*

assumes $I \neq \{\}$
assumes *finite* I
assumes *finite* E
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F ' I \subseteq \text{events}$
assumes $\text{prob} (\bigcap E) > 0$
obtains g **where** *bij-betw* $g \{0..<\text{card } I\} I$ **and** $\mathcal{P}((\bigcap (F ' g ' \{0..<\text{card } I\})) \mid (\bigcap E)) =$
 $(\prod i \in \{0..<\text{card } I\} . \mathcal{P}(F (g i) \mid (\bigcap ((F ' g ' \{0..<i\}) \cup E))))$
<proof>

lemma *prob-cond-Inter-index-cond-compl-fn:*

assumes $I \neq \{\}$
assumes *finite* I
assumes *finite* E
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F ' I \subseteq \text{events}$
assumes $\text{prob} (\bigcap E) > 0$
assumes *bb: bij-betw* $g \{0..<\text{card } I\} I$
shows $\mathcal{P}((\bigcap Aj \in I . \text{space } M - F Aj) \mid (\bigcap E)) =$
 $(\prod i \in \{0..<\text{card } I\} . \mathcal{P}(\text{space } M - F (g i) \mid (\bigcap (((\lambda Aj . \text{space } M - F Aj) ' g ' \{0..<i\}) \cup E))))$
<proof>

lemma *prob-cond-Inter-index-cond-compl-obtains:*

assumes $I \neq \{\}$
assumes *finite* I
assumes *finite* E
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F \text{ ' } I \subseteq \text{events}$
assumes $\text{prob} (\bigcap E) > 0$
obtains g **where** *bij-betw* $g \{0..<\text{card } I\} I$ **and** $\mathcal{P}((\bigcap Aj \in I . \text{space } M - F Aj)$
 $| (\bigcap E)) =$
 $(\prod i \in \{0..<\text{card } I\} . \mathcal{P}(\text{space } M - F (g i) | (\bigcap ((\lambda Aj . \text{space } M - F Aj) \text{ ' } g \text{ '}$
 $\{0..<i\} \cup E))))$
<proof>

lemma *prob-cond-inter-index-fn2:*

assumes $F \text{ ' } S \subseteq \text{events}$
assumes *finite* S
assumes $\text{card } S > 0$
assumes *bij-betw* $g \{0..<\text{card } S\} S$
shows $\text{prob} (\bigcap (F \text{ ' } S)) = \text{prob} (F (g 0)) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(F (g i) |$
 $(\bigcap (F \text{ ' } g \text{ ' } \{0..<i\}))))$
<proof>

lemma *prob-cond-inter-index-fn:*

assumes $F \text{ ' } S \subseteq \text{events}$
assumes *finite* S
assumes $S \neq \{\}$
assumes *bij-betw* $g \{0..<\text{card } S\} S$
shows $\text{prob} (\bigcap (F \text{ ' } S)) = \text{prob} (F (g 0)) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(F (g i) |$
 $(\bigcap (F \text{ ' } g \text{ ' } \{0..<i\}))))$
<proof>

lemma *prob-cond-inter-index-obtain-fn:*

assumes $F \text{ ' } S \subseteq \text{events}$
assumes *finite* S
assumes $S \neq \{\}$
obtains g **where** *bij-betw* $g \{0..<\text{card } S\} S$ **and**
 $\text{prob} (\bigcap (F \text{ ' } S)) = \text{prob} (F (g 0)) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(F (g i) | (\bigcap (F \text{ '}$
 $g \text{ ' } \{0..<i\}))))$
<proof>

lemma *prob-cond-inter-index-fn-compl:*

assumes $S \neq \{\}$
assumes $F \text{ ' } S \subseteq \text{events}$
assumes *finite* S
assumes *bij-betw* $f \{0..<\text{card } S\} S$
shows $\text{prob} (\bigcap ((-) (\text{space } M) \text{ ' } F \text{ ' } S)) = \text{prob} (\text{space } M - F (f 0)) *$
 $(\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(\text{space } M - F (f i) | (\bigcap ((-) (\text{space } M) \text{ ' } f \text{ '}$

{0..<i})))
 <proof>

lemma *prob-cond-inter-index-obtain-fn-compl:*

assumes $S \neq \{\}$
assumes $F \text{ ' } S \subseteq \text{events}$
assumes *finite* S
obtains f **where** *bij-betw* $f \{0..<card\ S\} S$ **and**
 $prob (\bigcap ((-) (space\ M) \text{ ' } F \text{ ' } S)) = prob (space\ M - F (f\ 0)) * (\prod i \in \{1..<(card\ S)\} . \mathcal{P}(space\ M - F (f\ i) \mid (\bigcap ((-) (space\ M) \text{ ' } F \text{ ' } f \text{ ' } \{0..<i\}))))$
 <proof>

lemma *prob-cond-Inter-take:*

assumes $S \neq \{\}$
assumes $S \subseteq \text{events}$
assumes *finite* S
obtains xs **where** *set* $xs = S$ **and** *length* $xs = card\ S$ **and**
 $prob (\bigcap S) = prob (hd\ xs) * (\prod i = 1..<(length\ xs) . \mathcal{P}((xs\ !\ i) \mid (\bigcap (set\ (take\ i\ xs))))))$
 <proof>

lemma *prob-cond-Inter-set-bound:*

assumes $A \neq \{\}$
assumes $A \subseteq \text{events}$
assumes *finite* A
assumes $\bigwedge Ai . f\ Ai \geq 0 \wedge f\ Ai \leq 1$
assumes $\bigwedge Ai\ S. Ai \in A \implies S \subseteq A - \{Ai\} \implies S \neq \{\} \implies \mathcal{P}(Ai \mid (\bigcap S)) \geq f\ Ai$
assumes $\bigwedge Ai. Ai \in A \implies prob\ Ai \geq f\ Ai$
shows $prob (\bigcap A) \geq (\prod a' \in A . f\ a')$
 <proof>
 end

end

5 Independent Events

theory *Indep-Events* **imports** *Cond-Prob-Extensions*
begin

5.1 More bijection helpers

lemma *bij-betw-obtain-subsetr:*

assumes *bij-betw* $f\ A\ B$
assumes $A' \subseteq A$
obtains B' **where** $B' \subseteq B$ **and** $B' = f \text{ ' } A'$

<proof>

lemma *bij-betw-obtain-subset*:

assumes *bij-betw* $f A B$

assumes $B' \subseteq B$

obtains A' **where** $A' \subseteq A$ **and** $B' = f \text{ ` } A'$

<proof>

lemma *bij-betw-remove*: $\text{bij-betw } f A B \implies a \in A \implies \text{bij-betw } f (A - \{a\}) (B - \{f a\})$

<proof>

5.2 Independent Event Extensions

Extensions on both the *indep_event* definition and the *indep_events* definition

context *prob-space*

begin

lemma *indep-eventsD*: $\text{indep-events } A I \implies (A \text{ ` } I \subseteq \text{events}) \implies J \subseteq I \implies J \neq \{\} \implies \text{finite } J \implies$

$\text{prob } (\bigcap_{j \in J}. A j) = (\prod_{j \in J}. \text{prob } (A j))$

<proof>

lemma

assumes *indep*: *indep-event* $A B$

shows *indep-eventD-ev1*: $A \in \text{events}$

and *indep-eventD-ev2*: $B \in \text{events}$

<proof>

lemma *indep-eventD*:

assumes *ie*: *indep-event* $A B$

shows $\text{prob } (A \cap B) = \text{prob } (A) * \text{prob } (B)$

<proof>

lemma *indep-eventI[intro]*:

assumes *ev*: $A \in \text{events } B \in \text{events}$

and *indep*: $\text{prob } (A \cap B) = \text{prob } A * \text{prob } B$

shows *indep-event* $A B$

<proof>

Alternate set definition - when no possibility of duplicate objects

definition *indep-events-set* :: '*a set* $\text{set} \implies \text{bool}$ **where**

indep-events-set $E \equiv (E \subseteq \text{events} \wedge (\forall J. J \subseteq E \longrightarrow \text{finite } J \longrightarrow J \neq \{\} \longrightarrow \text{prob } (\bigcap J) = (\prod_{i \in J}. \text{prob } i)))$

lemma *indep-events-setI[intro]*: $E \subseteq \text{events} \implies (\bigwedge J. J \subseteq E \implies \text{finite } J \implies J \neq \{\} \implies$

$\text{prob } (\bigcap J) = (\prod_{i \in J}. \text{prob } i) \implies \text{indep-events-set } E$

<proof>

lemma *indep-events-subset:*

indep-events-set $E \longleftrightarrow (\forall J \subseteq E. \text{indep-events-set } J)$

<proof>

lemma *indep-events-subset2:*

indep-events-set $E \implies J \subseteq E \implies \text{indep-events-set } J$

<proof>

lemma *indep-events-set-events:* *indep-events-set* $E \implies (\bigwedge e. e \in E \implies e \in \text{events})$

<proof>

lemma *indep-events-set-events-ss:* *indep-events-set* $E \implies E \subseteq \text{events}$

<proof>

lemma *indep-events-set-probs:* *indep-events-set* $E \implies J \subseteq E \implies \text{finite } J \implies J \neq \{\}$

$\text{prob } (\bigcap J) = (\prod_{i \in J} \text{prob } i)$

<proof>

lemma *indep-events-set-prod-all:* *indep-events-set* $E \implies \text{finite } E \implies E \neq \{\} \implies$

$\text{prob } (\bigcap E) = \text{prod prob } E$

<proof>

lemma *indep-events-not-contain-compl:*

assumes *indep-events-set* E

assumes $A \in E$

assumes $\text{prob } A > 0 \text{ prob } A < 1$

shows $(\text{space } M - A) \notin E$ (**is** $?A' \notin E$)

<proof>

lemma *indep-events-contain-compl-prob01:*

assumes *indep-events-set* E

assumes $A \in E$

assumes $\text{space } M - A \in E$

shows $\text{prob } A = 0 \vee \text{prob } A = 1$

<proof>

lemma *indep-events-set-singleton:*

assumes $A \in \text{events}$

shows *indep-events-set* $\{A\}$

<proof>

lemma *indep-events-pairs:*

assumes *indep-events-set* S

assumes $A \in S \ B \in S \ A \neq B$

shows *indep-event* $A B$
<proof>

lemma *indep-events-inter-pairs*:
assumes *indep-events-set* S
assumes *finite* A *finite* B
assumes $A \neq \{\}$ $B \neq \{\}$
assumes $A \subseteq S$ $B \subseteq S$ $A \cap B = \{\}$
shows *indep-event* $(\bigcap A)$ $(\bigcap B)$
<proof>

lemma *indep-events-inter-single*:
assumes *indep-events-set* S
assumes *finite* B
assumes $B \neq \{\}$
assumes $A \in S$ $B \subseteq S$ $A \notin B$
shows *indep-event* A $(\bigcap B)$
<proof>

lemma *indep-events-set-prob1*:
assumes $A \in \text{events}$
assumes *prob* $A = 1$
assumes $A \notin S$
assumes *indep-events-set* S
shows *indep-events-set* $(S \cup \{A\})$
<proof>

lemma *indep-events-set-prob0*:
assumes $A \in \text{events}$
assumes *prob* $A = 0$
assumes $A \notin S$
assumes *indep-events-set* S
shows *indep-events-set* $(S \cup \{A\})$
<proof>

lemma *indep-event-commute*:
assumes *indep-event* $A B$
shows *indep-event* $B A$
<proof>

Showing complement operation maintains independence

lemma *indep-event-one-compl*:
assumes *indep-event* $A B$
shows *indep-event* A $(\text{space } M - B)$
<proof>

lemma *indep-event-one-compl-rev*:
assumes $B \in \text{events}$

assumes *indep-event* A (*space* $M - B$)
shows *indep-event* $A B$
 \langle *proof* \rangle

lemma *indep-event-double-compl*: *indep-event* $A B \implies$ *indep-event* (*space* $M - A$) (*space* $M - B$)
 \langle *proof* \rangle

lemma *indep-event-double-compl-rev*: $A \in$ *events* $\implies B \in$ *events* \implies
indep-event (*space* $M - A$) (*space* $M - B$) \implies *indep-event* $A B$
 \langle *proof* \rangle

lemma *indep-events-set-one-compl*:
assumes *indep-events-set* S
assumes $A \in S$
shows *indep-events-set* ($\{\text{space } M - A\} \cup (S - \{A\})$)
 \langle *proof* \rangle

lemma *indep-events-set-update-compl*:
assumes *indep-events-set* E
assumes $E = A \cup B$
assumes $A \cap B = \{\}$
assumes *finite* E
shows *indep-events-set* ($((-) \text{ space } M - A) \cup B$)
 \langle *proof* \rangle

lemma *indep-events-set-compl*:
assumes *indep-events-set* E
assumes *finite* E
shows *indep-events-set* ($(\lambda e. \text{space } M - e) \text{ ' } E$)
 \langle *proof* \rangle

lemma *indep-event-empty*:
assumes $A \in$ *events*
shows *indep-event* $A \{\}$
 \langle *proof* \rangle

lemma *indep-event-compl-inter*:
assumes *indep-event* $A C$
assumes $B \in$ *events*
assumes *indep-event* $A (B \cap C)$
shows *indep-event* $A ((\text{space } M - B) \cap C)$
 \langle *proof* \rangle

lemma *indep-events-index-subset*:
indep-events $F E \longleftrightarrow (\forall J \subseteq E. \text{indep-events } F J)$

<proof>

lemma *indep-events-index-subset2:*

indep-events F E \implies J \subseteq E \implies indep-events F J

<proof>

lemma *indep-events-events-ss: indep-events F E \implies F ' E \subseteq events*

<proof>

lemma *indep-events-events: indep-events F E \implies ($\bigwedge e. e \in E \implies F e \in \text{events}$)*

<proof>

lemma *indep-events-probs: indep-events F E \implies J \subseteq E \implies finite J \implies J \neq {} \implies prob ($\bigcap (F ' J)$) = ($\prod_{i \in J. \text{prob } (F i)}$)*

<proof>

lemma *indep-events-prod-all: indep-events F E \implies finite E \implies E \neq {} \implies prob ($\bigcap (F ' E)$) = ($\prod_{i \in E. \text{prob } (F i)}$)*

<proof>

lemma *indep-events-ev-not-contain-compl:*

assumes *indep-events F E*

assumes *A \in E*

assumes *prob (F A) > 0 prob (F A) < 1*

shows *(space M - F A) \notin F ' E (is ?A' \notin F ' E)*

<proof>

lemma *indep-events-singleton:*

assumes *F A \in events*

shows *indep-events F {A}*

<proof>

lemma *indep-events-ev-pairs:*

assumes *indep-events F S*

assumes *A \in S B \in S A \neq B*

shows *indep-event (F A) (F B)*

<proof>

lemma *indep-events-ev-inter-pairs:*

assumes *indep-events F S*

assumes *finite A finite B*

assumes *A \neq {} B \neq {}*

assumes *A \subseteq S B \subseteq S A \cap B = {}*

shows *indep-event ($\bigcap (F ' A)$) ($\bigcap (F ' B)$)*

<proof>

lemma *indep-events-ev-inter-single:*

assumes *indep-events* $F S$
assumes *finite* B
assumes $B \neq \{\}$
assumes $A \in S \ B \subseteq S \ A \notin B$
shows *indep-event* $(F A) (\bigcap (F \text{ ` } B))$
 <proof>

lemma *indep-events-fn-eq*:
assumes $\bigwedge Ai. Ai \in E \implies F Ai = G Ai$
assumes *indep-events* $F E$
shows *indep-events* $G E$
 <proof>

lemma *indep-events-fn-eq-iff*:
assumes $\bigwedge Ai. Ai \in E \implies F Ai = G Ai$
shows *indep-events* $F E \longleftrightarrow \text{indep-events } G E$
 <proof>

lemma *indep-events-one-compl*:
assumes *indep-events* $F S$
assumes $A \in S$
shows *indep-events* $(\lambda i. \text{if } (i = A) \text{ then } (\text{space } M - F i) \text{ else } F i) S$ (**is**
indep-events $?G S$)
 <proof>

lemma *indep-events-update-compl*:
assumes *indep-events* $F E$
assumes $E = A \cup B$
assumes $A \cap B = \{\}$
assumes *finite* E
shows *indep-events* $(\lambda Ai. \text{if } (Ai \in A) \text{ then } (\text{space } M - (F Ai)) \text{ else } (F Ai)) E$
 <proof>

lemma *indep-events-compl*:
assumes *indep-events* $F E$
assumes *finite* E
shows *indep-events* $(\lambda Ai. \text{space } M - F Ai) E$
 <proof>

lemma *indep-events-impl-inj-on*:
assumes *finite* A
assumes *indep-events* $F A$
assumes $\bigwedge A'. A' \in A \implies \text{prob } (F A') > 0 \wedge \text{prob } (F A') < 1$
shows *inj-on* $F A$
 <proof>

lemma *indep-events-imp-set*:
assumes *finite* A
assumes *indep-events* $F A$

assumes $\bigwedge A' . A' \in A \implies \text{prob } (F A') > 0 \wedge \text{prob } (F A') < 1$
shows *indep-events-set* $(F A)$
 <proof>

lemma *indep-event-set-equiv-bij*:
assumes *bij-betw* $F A E$
assumes *finite* E
shows *indep-events-set* $E \longleftrightarrow \text{indep-events } F A$
 <proof>

5.3 Mutual Independent Events

Note, set based version only if no duplicates in usage case. The `mutual_indep_events` definition is more general and recommended

definition *mutual-indep-set*:: 'a set \implies 'a set set \implies bool
where *mutual-indep-set* $A S \longleftrightarrow A \in \text{events} \wedge S \subseteq \text{events} \wedge (\forall T \subseteq S . T \neq \{\} \longrightarrow \text{prob } (A \cap (\bigcap T)) = \text{prob } A * \text{prob } (\bigcap T))$

lemma *mutual-indep-setI[intro]*: $A \in \text{events} \implies S \subseteq \text{events} \implies (\bigwedge T . T \subseteq S \implies T \neq \{\}) \implies \text{prob } (A \cap (\bigcap T)) = \text{prob } A * \text{prob } (\bigcap T) \implies \text{mutual-indep-set } A S$
 <proof>

lemma *mutual-indep-setD[dest]*: $\text{mutual-indep-set } A S \implies T \subseteq S \implies T \neq \{\} \implies \text{prob } (A \cap (\bigcap T)) = \text{prob } A * \text{prob } (\bigcap T)$
 <proof>

lemma *mutual-indep-setD2[dest]*: $\text{mutual-indep-set } A S \implies A \in \text{events}$
 <proof>

lemma *mutual-indep-setD3[dest]*: $\text{mutual-indep-set } A S \implies S \subseteq \text{events}$
 <proof>

lemma *mutual-indep-subset*: $\text{mutual-indep-set } A S \implies T \subseteq S \implies \text{mutual-indep-set } A T$
 <proof>

lemma *mutual-indep-event-set-defD*:
assumes *mutual-indep-set* $A S$
assumes *finite* T
assumes $T \subseteq S$
assumes $T \neq \{\}$
shows *indep-event* $A (\bigcap T)$
 <proof>

lemma *mutual-indep-event-defI*: $A \in \text{events} \implies S \subseteq \text{events} \implies (\bigwedge T . T \subseteq S \implies T \neq \{\}) \implies$

indep-event $A (\bigcap T) \implies$ *mutual-indep-set* $A S$
<proof>

lemma *mutual-indep-singleton-event*: *mutual-indep-set* $A S \implies B \in S \implies$ *indep-event* $A B$
<proof>

lemma *mutual-indep-cond*:
assumes $A \in \text{events}$ **and** $T \subseteq \text{events}$ **and** *finite* T
and *mutual-indep-set* $A S$ **and** $T \subseteq S$ **and** $T \neq \{\}$ **and** $\text{prob} (\bigcap T) \neq 0$
shows $\mathcal{P}(A | (\bigcap T)) = \text{prob } A$
<proof>

lemma *mutual-indep-cond-full*:
assumes $A \in \text{events}$ **and** $S \subseteq \text{events}$ **and** *finite* S
and *mutual-indep-set* $A S$ **and** $S \neq \{\}$ **and** $\text{prob} (\bigcap S) \neq 0$
shows $\mathcal{P}(A | (\bigcap S)) = \text{prob } A$
<proof>

lemma *mutual-indep-cond-single*:
assumes $A \in \text{events}$ **and** $B \in \text{events}$
and *mutual-indep-set* $A S$ **and** $B \in S$ **and** $\text{prob } B \neq 0$
shows $\mathcal{P}(A | B) = \text{prob } A$
<proof>

lemma *mutual-indep-set-empty*: $A \in \text{events} \implies$ *mutual-indep-set* $A \{\}$
<proof>

lemma *not-mutual-indep-set-itself*:
assumes $\text{prob } A > 0$ **and** $\text{prob } A < 1$
shows \neg *mutual-indep-set* $A \{A\}$
<proof>

lemma *is-mutual-indep-set-itself*:
assumes $A \in \text{events}$
assumes $\text{prob } A = 0 \vee \text{prob } A = 1$
shows *mutual-indep-set* $A \{A\}$
<proof>

lemma *mutual-indep-set-singleton*:
assumes *indep-event* $A B$
shows *mutual-indep-set* $A \{B\}$
<proof>

lemma *mutual-indep-set-one-compl*:
assumes *mutual-indep-set* $A S$
assumes *finite* S
assumes $B \in S$
shows *mutual-indep-set* $A (\{\text{space } M - B\} \cup S)$

<proof>

lemma *mutual-indep-events-set-update-compl:*

assumes *mutual-indep-set* $X E$

assumes $E = A \cup B$

assumes $A \cap B = \{\}$

assumes *finite* E

shows *mutual-indep-set* $X ((-) (space M) ' A) \cup B$

<proof>

lemma *mutual-indep-events-compl:*

assumes *finite* S

assumes *mutual-indep-set* $A S$

shows *mutual-indep-set* $A ((\lambda s . space M - s) ' S)$

<proof>

lemma *mutual-indep-set-all:*

assumes $A \subseteq events$

assumes $\bigwedge Ai. Ai \in A \implies (mutual-indep-set Ai (A - \{Ai\}))$

shows *indep-events-set* A

<proof>

Prefered version using indexed notation

definition *mutual-indep-events:: 'a set \implies (nat \implies 'a set) \implies nat set \implies bool*

where *mutual-indep-events* $A F I \longleftrightarrow A \in events \wedge (F ' I \subseteq events) \wedge (\forall J \subseteq I . J \neq \{\} \implies prob (A \cap (\bigcap j \in J . F j)) = prob A * prob (\bigcap j \in J . F j))$

lemma *mutual-indep-eventsI[intro]:* $A \in events \implies (F ' I \subseteq events) \implies (\bigwedge J. J \subseteq I \implies J \neq \{\} \implies$

$prob (A \cap (\bigcap j \in J . F j)) = prob A * prob (\bigcap j \in J . F j)) \implies mutual-indep-events A F I$

<proof>

lemma *mutual-indep-eventsD[dest]:* $mutual-indep-events A F I \implies J \subseteq I \implies J \neq \{\} \implies prob (A \cap (\bigcap j \in J . F j)) = prob A * prob (\bigcap j \in J . F j)$

<proof>

lemma *mutual-indep-eventsD2[dest]:* $mutual-indep-events A F I \implies A \in events$

<proof>

lemma *mutual-indep-eventsD3[dest]:* $mutual-indep-events A F I \implies F ' I \subseteq events$

<proof>

lemma *mutual-indep-ev-subset:* $mutual-indep-events A F I \implies J \subseteq I \implies mutual-indep-events A F J$

<proof>

lemma *mutual-indep-event-defD*:
assumes *mutual-indep-events A F I*
assumes *finite J*
assumes $J \subseteq I$
assumes $J \neq \{\}$
shows *indep-event A* $(\bigcap_{j \in J} . F j)$
 \langle *proof* \rangle

lemma *mutual-ev-indep-event-defI*: $A \in \text{events} \implies F \text{ ' } I \subseteq \text{events} \implies (\bigwedge J. J \subseteq I \implies J \neq \{\}) \implies$
indep-event A $(\bigcap (F \text{ ' } J)) \implies$ *mutual-indep-events A F I*
 \langle *proof* \rangle

lemma *mutual-indep-ev-singleton-event*:
assumes *mutual-indep-events A F I*
assumes $B \in F \text{ ' } I$
shows *indep-event A B*
 \langle *proof* \rangle

lemma *mutual-indep-ev-singleton-event2*:
assumes *mutual-indep-events A F I*
assumes $i \in I$
shows *indep-event A* $(F i)$
 \langle *proof* \rangle

lemma *mutual-indep-iff*:
shows *mutual-indep-events A F I* \longleftrightarrow *mutual-indep-set A* $(F \text{ ' } I)$
 \langle *proof* \rangle

lemma *mutual-indep-ev-cond*:
assumes $A \in \text{events}$ **and** $F \text{ ' } J \subseteq \text{events}$ **and** *finite J*
and *mutual-indep-events A F I* **and** $J \subseteq I$ **and** $J \neq \{\}$ **and** $\text{prob} (\bigcap (F \text{ ' } J)) \neq 0$
shows $\mathcal{P}(A \mid (\bigcap (F \text{ ' } J))) = \text{prob } A$
 \langle *proof* \rangle

lemma *mutual-indep-ev-cond-full*:
assumes $A \in \text{events}$ **and** $F \text{ ' } I \subseteq \text{events}$ **and** *finite I*
and *mutual-indep-events A F I* **and** $I \neq \{\}$ **and** $\text{prob} (\bigcap (F \text{ ' } I)) \neq 0$
shows $\mathcal{P}(A \mid (\bigcap (F \text{ ' } I))) = \text{prob } A$
 \langle *proof* \rangle

lemma *mutual-indep-ev-cond-single*:
assumes $A \in \text{events}$ **and** $B \in \text{events}$
and *mutual-indep-events A F I* **and** $B \in F \text{ ' } I$ **and** $\text{prob } B \neq 0$
shows $\mathcal{P}(A \mid B) = \text{prob } A$
 \langle *proof* \rangle

lemma *mutual-indep-ev-empty*: $A \in \text{events} \implies$ *mutual-indep-events A F* $\{\}$
 \langle *proof* \rangle

lemma *not-mutual-indep-ev-itself*:

assumes $\text{prob } A > 0$ **and** $\text{prob } A < 1$ **and** $A = F\ i$

shows $\neg \text{mutual-indep-events } A\ F\ \{i\}$

<proof>

lemma *is-mutual-indep-ev-itself*:

assumes $A \in \text{events}$ **and** $A = F\ i$

assumes $\text{prob } A = 0 \vee \text{prob } A = 1$

shows $\text{mutual-indep-events } A\ F\ \{i\}$

<proof>

lemma *mutual-indep-ev-singleton*:

assumes $\text{indep-event } A\ (F\ i)$

shows $\text{mutual-indep-events } A\ F\ \{i\}$

<proof>

lemma *mutual-indep-ev-one-compl*:

assumes $\text{mutual-indep-events } A\ F\ I$

assumes $\text{finite } I$

assumes $i \in I$

assumes $\text{space } M - F\ i = F\ j$

shows $\text{mutual-indep-events } A\ F\ (\{j\} \cup I)$

<proof>

lemma *mutual-indep-events-update-compl*:

assumes $\text{mutual-indep-events } X\ F\ S$

assumes $S = A \cup B$

assumes $A \cap B = \{\}$

assumes $\text{finite } S$

assumes $\text{bij-betw } G\ A\ A'$

assumes $\bigwedge i. i \in A \implies F\ (G\ i) = \text{space } M - F\ i$

shows $\text{mutual-indep-events } X\ F\ (A' \cup B)$

<proof>

lemma *mutual-indep-ev-events-compl*:

assumes $\text{finite } S$

assumes $\text{mutual-indep-events } A\ F\ S$

assumes $\text{bij-betw } G\ S\ S'$

assumes $\bigwedge i. i \in S \implies F\ (G\ i) = \text{space } M - F\ i$

shows $\text{mutual-indep-events } A\ F\ S'$

<proof>

Important lemma on relation between independence and mutual independence of a set

lemma *mutual-indep-ev-set-all*:

assumes $F\ 'I \subseteq \text{events}$

assumes $\bigwedge i. i \in I \implies (\text{mutual-indep-events } (F\ i)\ F\ (I - \{i\}))$

shows $\text{indep-events } F\ I$

<proof>

end
end

6 The Basic Probabilistic Method Framework

This theory includes all aspects of step (3) and (4) of the basic method framework, which are purely probabilistic

theory *Basic-Method* **imports** *Indep-Events*
begin

6.1 More Set and Multiset lemmas

lemma *card-size-set-mset*: $\text{card } (\text{set-mset } A) \leq \text{size } A$

<proof>

lemma *Union-exists*: $\{a \in A . \exists b \in B . P a b\} = (\bigcup b \in B . \{a \in A . P a b\})$

<proof>

lemma *Inter-forall*: $B \neq \{\} \implies \{a \in A . \forall b \in B . P a b\} = (\bigcap b \in B . \{a \in A . P a b\})$

<proof>

lemma *function-map-multi-filter-size*:

assumes *image-mset* F (*mset-set* A) = B **and** *finite* A

shows $\text{card } \{a \in A . P (F a)\} = \text{size } \{\# b \in \# B . P b \#\}$

<proof>

lemma *bij-mset-obtain-set-elem*:

assumes *image-mset* F (*mset-set* A) = B

assumes $b \in \# B$

obtains a **where** $a \in A$ **and** $F a = b$

<proof>

lemma *bij-mset-obtain-mset-elem*:

assumes *finite* A

assumes *image-mset* F (*mset-set* A) = B

assumes $a \in A$

obtains b **where** $b \in \# B$ **and** $F a = b$

<proof>

lemma *prod-fn-le1*:

fixes $f :: 'c \Rightarrow ('d :: \{\text{comm-monoid-mult, linordered-semidom}\})$

assumes *finite* A

assumes $A \neq \{\}$

assumes $\bigwedge y. y \in A \implies f y \geq 0 \wedge f y < 1$

shows $(\prod x \in A. f x) < 1$

<proof>

context *prob-space*
begin

6.2 Existence Lemmas

lemma *prob-lt-one-obtain:*

assumes $\{e \in \text{space } M . Q\ e\} \in \text{events}$
assumes $\text{prob } \{e \in \text{space } M . Q\ e\} < 1$
obtains e **where** $e \in \text{space } M$ **and** $\neg Q\ e$

<proof>

lemma *prob-gt-zero-obtain:*

assumes $\{e \in \text{space } M . Q\ e\} \in \text{events}$
assumes $\text{prob } \{e \in \text{space } M . Q\ e\} > 0$
obtains e **where** $e \in \text{space } M$ **and** $Q\ e$

<proof>

lemma *inter-gt0-event:*

assumes $F\ 'I \subseteq \text{events}$
assumes $\text{prob } (\bigcap i \in I . (\text{space } M - (F\ i))) > 0$
shows $(\bigcap i \in I . (\text{space } M - (F\ i))) \in \text{events}$ **and** $(\bigcap i \in I . (\text{space } M - (F\ i))) \neq \{\}$

<proof>

lemma *obtain-intersection:*

assumes $F\ 'I \subseteq \text{events}$
assumes $\text{prob } (\bigcap i \in I . (\text{space } M - (F\ i))) > 0$
obtains e **where** $e \in \text{space } M$ **and** $\bigwedge i . i \in I \implies e \notin F\ i$

<proof>

lemma *obtain-intersection-prop:*

assumes $F\ 'I \subseteq \text{events}$
assumes $\bigwedge i . i \in I \implies F\ i = \{e \in \text{space } M . P\ e\ i\}$
assumes $\text{prob } (\bigcap i \in I . (\text{space } M - (F\ i))) > 0$
obtains e **where** $e \in \text{space } M$ **and** $\bigwedge i . i \in I \implies \neg P\ e\ i$

<proof>

lemma *not-in-big-union:*

assumes $\bigwedge i . i \in A \implies e \notin i$
shows $e \notin (\bigcup A)$

<proof>

lemma *not-in-big-union-fn:*

assumes $\bigwedge i . i \in A \implies e \notin F\ i$
shows $e \notin (\bigcup i \in A . F\ i)$

<proof>

lemma *obtain-intersection-union:*

assumes $F' I \subseteq \text{events}$

assumes $\text{prob} (\bigcap_{i \in I} (\text{space } M - (F' i))) > 0$

obtains e **where** $e \in \text{space } M$ **and** $e \notin (\bigcup_{i \in I} F' i)$

<proof>

6.3 Basic Bounds

Lemmas on the Complete Independence and Union bound

lemma *complete-indep-bound1:*

assumes *finite* A

assumes $A \neq \{\}$

assumes $A \subseteq \text{events}$

assumes *indep-events-set* A

assumes $\bigwedge a . a \in A \implies \text{prob } a < 1$

shows $\text{prob} (\text{space } M - (\bigcap A)) > 0$

<proof>

lemma *complete-indep-bound1-index:*

assumes *finite* A

assumes $A \neq \{\}$

assumes $F' A \subseteq \text{events}$

assumes *indep-events* $F A$

assumes $\bigwedge a . a \in A \implies \text{prob} (F a) < 1$

shows $\text{prob} (\text{space } M - (\bigcap (F' A))) > 0$

<proof>

lemma *complete-indep-bound2:*

assumes *finite* A

assumes $A \subseteq \text{events}$

assumes *indep-events-set* A

assumes $\bigwedge a . a \in A \implies \text{prob } a < 1$

shows $\text{prob} (\text{space } M - (\bigcup A)) > 0$

<proof>

lemma *complete-indep-bound2-index:*

assumes *finite* A

assumes $F' A \subseteq \text{events}$

assumes *indep-events* $F A$

assumes $\bigwedge a . a \in A \implies \text{prob} (F a) < 1$

shows $\text{prob} (\text{space } M - (\bigcup (F' A))) > 0$

<proof>

lemma *complete-indep-bound3:*

assumes *finite* A

assumes $A \neq \{\}$

assumes $F' A \subseteq \text{events}$

assumes *indep-events* $F A$

assumes $\bigwedge a . a \in A \implies \text{prob} (F a) < 1$

shows $\text{prob} (\bigcap a \in A. \text{space } M - F a) > 0$
<proof>

Combining complete independence with existence step

lemma *complete-indep-bound-obtain:*
assumes *finite A*
assumes $A \subseteq \text{events}$
assumes *indep-events-set A*
assumes $\bigwedge a . a \in A \implies \text{prob } a < 1$
obtains *e* **where** $e \in \text{space } M$ **and** $e \notin \bigcup A$
<proof>

lemma *Union-bound-events:*
assumes *finite A*
assumes $A \subseteq \text{events}$
shows $\text{prob} (\bigcup A) \leq (\sum a \in A. \text{prob } a)$
<proof>

lemma *Union-bound-events-fun:*
assumes *finite A*
assumes $f ' A \subseteq \text{events}$
shows $\text{prob} (\bigcup (f ' A)) \leq (\sum a \in A. \text{prob } (f a))$
<proof>

lemma *Union-bound-avoid:*
assumes *finite A*
assumes $(\sum a \in A. \text{prob } a) < 1$
assumes $A \subseteq \text{events}$
shows $\text{prob} (\text{space } M - \bigcup A) > 0$
<proof>

lemma *Union-bound-avoid-fun:*
assumes *finite A*
assumes $(\sum a \in A. \text{prob } (f a)) < 1$
assumes $f ' A \subseteq \text{events}$
shows $\text{prob} (\text{space } M - \bigcup (f ' A)) > 0$
<proof>

Combining union bound with existence step

lemma *Union-bound-obtain:*
assumes *finite A*
assumes $(\sum a \in A. \text{prob } a) < 1$
assumes $A \subseteq \text{events}$
obtains *e* **where** $e \in \text{space } M$ **and** $e \notin \bigcup A$
<proof>

lemma *Union-bound-obtain-fun:*
assumes *finite A*

assumes $(\sum a \in A. \text{prob } (f a)) < 1$
assumes $f' A \subseteq \text{events}$
obtains e **where** $e \in \text{space } M$ **and** $e \notin \bigcup (f' A)$
 <proof>

lemma *Union-bound-obtain-compl:*
assumes *finite* A
assumes $(\sum a \in A. \text{prob } a) < 1$
assumes $A \subseteq \text{events}$
obtains e **where** $e \in (\text{space } M - \bigcup A)$
 <proof>

lemma *Union-bound-obtain-compl-fun:*
assumes *finite* A
assumes $(\sum a \in A. \text{prob } (f a)) < 1$
assumes $f' A \subseteq \text{events}$
obtains e **where** $e \in (\text{space } M - \bigcup (f' A))$
 <proof>

end

end

7 Lovasz Local Lemma

theory *Lovasz-Local-Lemma*
imports
 Basic-Method
 HOL-Real-Asymp.Real-Asymp
 Indep-Events
 Digraph-Extensions
begin

7.1 Random Lemmas on Product Operator

lemma *prod-constant-ge:*
fixes $y :: 'b :: \{\text{comm-monoid-mult}, \text{linordered-semidom}\}$
assumes $\text{card } A \leq k$
assumes $y \geq 0$ **and** $y < 1$
shows $(\prod x \in A. y) \geq y \wedge k$
 <proof>

lemma **(in** *linordered-idom*) *prod-mono3:*
assumes *finite* J $I \subseteq J$ $\bigwedge i. i \in J \implies 0 \leq f i$ $(\bigwedge i. i \in J \implies f i \leq 1)$
shows $\text{prod } f J \leq \text{prod } f I$
 <proof>

lemma *bij-on-ss-image:*
assumes $A \subseteq B$

assumes *bij-betw* g B B'
shows $g \text{ ' } A \subseteq B'$
 ⟨*proof*⟩

lemma *bij-on-ss-proper-image*:

assumes $A \subset B$
assumes *bij-betw* g B B'
shows $g \text{ ' } A \subset B'$
 ⟨*proof*⟩

7.2 Dependency Graph Concept

Uses directed graphs. The `pair_digraph` locale was sufficient as multi-edges are irrelevant

locale *dependency-digraph* = *pair-digraph* G :: *nat pair-pre-digraph* + *prob-space*
 M :: 'a *measure*

for G M + **fixes** F :: *nat* \Rightarrow 'a *set*
assumes *uss*: $F \text{ ' } (pverts\ G) \subseteq events$
assumes *mis*: $\bigwedge i. i \in (pverts\ G) \implies mutual-indep-events\ (F\ i)\ F\ ((pverts\ G) - (\{i\} \cup neighborhood\ i))$
begin

lemma *dep-graph-indiv-nh-indep*:

assumes $A \in pverts\ G$ $B \in pverts\ G$
assumes $B \notin neighborhood\ A$
assumes $A \neq B$
assumes *prob* $(F\ B) \neq 0$
shows $\mathcal{P}((F\ A) \mid (F\ B)) = prob\ (F\ A)$
 ⟨*proof*⟩

lemma *mis-subset*:

assumes $i \in pverts\ G$
assumes $A \subseteq pverts\ G$
shows *mutual-indep-events* $(F\ i)\ F\ (A - (\{i\} \cup neighborhood\ i))$
 ⟨*proof*⟩

lemma *dep-graph-indep-events*:

assumes $A \subseteq pverts\ G$
assumes $\bigwedge Ai. Ai \in A \implies out-degree\ G\ Ai = 0$
shows *indep-events* $F\ A$
 ⟨*proof*⟩

end

7.3 Lovasz Local General Lemma

context *prob-space*

begin

lemma compl-sets-index:

assumes $F \text{ ' } A \subseteq \text{events}$

shows $(\lambda i. \text{space } M - F i) \text{ ' } A \subseteq \text{events}$

$\langle \text{proof} \rangle$

lemma lovasz-inductive-base:

assumes $\text{dependency-digraph } G M F$

assumes $\bigwedge Ai. Ai \in A \implies g Ai \geq 0 \wedge g Ai < 1$

assumes $\bigwedge Ai. Ai \in A \implies (\text{prob } (F Ai) \leq (g Ai) * (\prod_{Aj \in \text{pre-digraph.neighborhood } G Ai. (1 - (g Aj))}))$

assumes $Ai \in A$

assumes $\text{pverts } G = A$

shows $\text{prob } (F Ai) \leq g Ai$

$\langle \text{proof} \rangle$

lemma lovasz-inductive-base-set:

assumes $N \subseteq A$

assumes $\bigwedge Ai. Ai \in A \implies g Ai \geq 0 \wedge g Ai < 1$

assumes $\bigwedge Ai. Ai \in A \implies (\text{prob } (F Ai) \leq (g Ai) * (\prod_{Aj \in N. (1 - (g Aj))}))$

assumes $Ai \in A$

shows $\text{prob } (F Ai) \leq g Ai$

$\langle \text{proof} \rangle$

lemma split-prob-lt-helper:

assumes $\text{dep-graph: dependency-digraph } G M F$

assumes $\text{dep-graph-verts: pverts } G = A$

assumes $\text{fbounds: } \bigwedge i. i \in A \implies f i \geq 0 \wedge f i < 1$

assumes $\text{prob-Ai: } \bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq$

$(f Ai) * (\prod_{Aj \in \text{pre-digraph.neighborhood } G Ai. (1 - (f Aj))})$

assumes $\text{aiin: } Ai \in A$

assumes $N \subseteq \text{pre-digraph.neighborhood } G Ai$

assumes $\exists P1 P2. \mathcal{P}(F Ai \mid \bigcap_{Aj \in S. \text{space } M - F Aj}) = P1/P2 \wedge$

$P1 \leq \text{prob } (F Ai) \wedge P2 \geq (\prod_{Aj \in N. (1 - (f Aj))})$

shows $\mathcal{P}(F Ai \mid \bigcap_{Aj \in S. \text{space } M - F Aj}) \leq f Ai$

$\langle \text{proof} \rangle$

lemma lovasz-inequality:

assumes $\text{finS: finite } S$

assumes $\text{sevents: } F \text{ ' } S \subseteq \text{events}$

assumes $\text{S-subset: } S \subseteq A - \{Ai\}$

assumes $\text{prob2: } \text{prob } (\bigcap_{Aj \in S. (\text{space } M - (F Aj))) > 0$

assumes $\text{irange: } i \in \{0..<\text{card } S1\}$

assumes $\text{bb: bij-betw } g \{0..<\text{card } S1\} S1$

assumes $\text{s1-def: } S1 = (S \cap N)$

assumes $\text{s2-def: } S2 = S - S1$

assumes $\text{ne-cond: } i > 0 \vee S2 \neq \{\}$

assumes $\text{hyps: } \bigwedge B. B \subset S \implies g i \in A \implies B \subseteq A - \{g i\} \implies B \neq \{\} \implies$

$0 < \text{prob } (\bigcap_{Aj \in B. \text{space } M - F Aj}) \implies \mathcal{P}(F (g i) \mid \bigcap_{Aj \in B. \text{space } M - F Aj}) \leq f (g i)$

shows $\mathcal{P}((\text{space } M - F(g\ i)) \mid (\bigcap ((\lambda\ i.\ \text{space } M - F\ i) \text{ ' } g \text{ ' } \{0..<i\} \cup ((\lambda\ i.\ \text{space } M - F\ i) \text{ ' } S2))))$
 $\geq (1 - f(g\ i))$
 ⟨proof⟩

The main helper lemma

lemma *lovasz-inductive*:

assumes *finA*: *finite A*
assumes *Aevents*: $F \text{ ' } A \subseteq \text{events}$
assumes *fbounds*: $\bigwedge i.\ i \in A \implies f\ i \geq 0 \wedge f\ i < 1$
assumes *dep-graph*: *dependency-digraph G M F*
assumes *dep-graph-verts*: *pverts G = A*
assumes *prob-Ai*: $\bigwedge Ai.\ Ai \in A \implies \text{prob}(F\ Ai) \leq$
 $(f\ Ai) * (\prod Aj \in \text{pre-digraph.neighborhood } G\ Ai.\ (1 - (f\ Aj)))$
assumes *Ai-in*: $Ai \in A$
assumes *S-subset*: $S \subseteq A - \{Ai\}$
assumes *S-nempty*: $S \neq \{\}$
assumes *prob2*: $\text{prob}(\bigcap Aj \in S.\ (\text{space } M - (F\ Aj))) > 0$
shows $\mathcal{P}((F\ Ai) \mid (\bigcap Aj \in S.\ (\text{space } M - (F\ Aj)))) \leq f\ Ai$
 ⟨proof⟩

The main lemma

theorem *lovasz-local-general*:

assumes $A \neq \{\}$
assumes $F \text{ ' } A \subseteq \text{events}$
assumes *finite A*
assumes $\bigwedge Ai.\ Ai \in A \implies f\ Ai \geq 0 \wedge f\ Ai < 1$
assumes *dependency-digraph G M F*
assumes $\bigwedge Ai.\ Ai \in A \implies (\text{prob}(F\ Ai) \leq (f\ Ai) * (\prod Aj \in \text{pre-digraph.neighborhood } G\ Ai.\ (1 - (f\ Aj))))$
assumes *pverts G = A*
shows $\text{prob}(\bigcap Ai \in A.\ (\text{space } M - (F\ Ai))) \geq (\prod Ai \in A.\ (1 - f\ Ai)) (\prod Ai \in A.\ (1 - f\ Ai)) > 0$
 ⟨proof⟩

7.4 Lovasz Corollaries and Variations

corollary *lovasz-local-general-positive*:

assumes $A \neq \{\}$
assumes $F \text{ ' } A \subseteq \text{events}$
assumes *finite A*
assumes $\bigwedge Ai.\ Ai \in A \implies f\ Ai \geq 0 \wedge f\ Ai < 1$
assumes *dependency-digraph G M F*
assumes $\bigwedge Ai.\ Ai \in A \implies (\text{prob}(F\ Ai) \leq$
 $(f\ Ai) * (\prod Aj \in \text{pre-digraph.neighborhood } G\ Ai.\ (1 - (f\ Aj))))$
assumes *pverts G = A*
shows $\text{prob}(\bigcap Ai \in A.\ (\text{space } M - (F\ Ai))) > 0$
 ⟨proof⟩

theorem *lovasz-local-symmetric-dep-graph*:
fixes $e :: \text{real}$
fixes $d :: \text{nat}$
assumes $A \neq \{\}$
assumes $F \text{ ' } A \subseteq \text{events}$
assumes *finite* A
assumes *dependency-digraph* $G M F$
assumes $\bigwedge Ai. Ai \in A \implies \text{out-degree } G Ai \leq d$
assumes $\bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq p$
assumes $\exp(1) * p * (d + 1) \leq 1$
assumes *pverts* $G = A$
shows $\text{prob } (\bigcap Ai \in A . (\text{space } M - (F Ai))) > 0$
<proof>

corollary *lovasz-local-symmetric4gt*:
fixes $e :: \text{real}$
fixes $d :: \text{nat}$
assumes $A \neq \{\}$
assumes $F \text{ ' } A \subseteq \text{events}$
assumes *finite* A
assumes *dependency-digraph* $G M F$
assumes $\bigwedge Ai. Ai \in A \implies \text{out-degree } G Ai \leq d$
assumes $\bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq p$
assumes $\frac{1}{4} * p * d \leq 1$
assumes $d \geq 3$
assumes *pverts* $G = A$
shows $\text{prob } (\bigcap Ai \in A . (\text{space } M - F Ai)) > 0$
<proof>

lemma *lovasz-local-symmetric4*:
fixes $e :: \text{real}$
fixes $d :: \text{nat}$
assumes $A \neq \{\}$
assumes $F \text{ ' } A \subseteq \text{events}$
assumes *finite* A
assumes *dependency-digraph* $G M F$
assumes $\bigwedge Ai. Ai \in A \implies \text{out-degree } G Ai \leq d$
assumes $\bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq p$
assumes $\frac{1}{4} * p * d \leq 1$
assumes $d \geq 1$
assumes *pverts* $G = A$
shows $\text{prob } (\bigcap Ai \in A . (\text{space } M - F Ai)) > 0$
<proof>

Converting between dependency graph and indexed set representation of mutual independence

lemma (*in pair-digraph*) *g-Ai-simplification*:
assumes $Ai \in A$

assumes $g\ Ai \subseteq A - \{Ai\}$
assumes $pverts\ G = A$
assumes $parcs\ G = \{e \in A \times A . snd\ e \in (A - (\{fst\ e\} \cup (g\ (fst\ e))))\}$
shows $g\ Ai = A - (\{Ai\} \cup neighborhood\ Ai)$
 <proof>

lemma *define-dep-graph-set:*

assumes $A \neq \{\}$
assumes $F\ 'A \subseteq events$
assumes *finite* A
assumes $\bigwedge Ai. Ai \in A \implies g\ Ai \subseteq A - \{Ai\} \wedge mutual-indep-events\ (F\ Ai)\ F\ (g\ Ai)$
shows $dependency-digraph\ (\big\| pverts = A, parcs = \{e \in A \times A . snd\ e \in (A - (\{fst\ e\} \cup (g\ (fst\ e))))\} \big\|)\ M\ F$
 (is *dependency-digraph* $?G\ M\ F$)
 <proof>

lemma *define-dep-graph-deg-bound:*

assumes $A \neq \{\}$
assumes $F\ 'A \subseteq events$
assumes *finite* A
assumes $\bigwedge Ai. Ai \in A \implies g\ Ai \subseteq A - \{Ai\} \wedge card\ (g\ Ai) \geq card\ A - d - 1$
 \wedge
 $mutual-indep-events\ (F\ Ai)\ F\ (g\ Ai)$
shows $\bigwedge Ai. Ai \in A \implies$
 $out-degree\ (\big\| pverts = A, parcs = \{e \in A \times A . snd\ e \in (A - (\{fst\ e\} \cup (g\ (fst\ e))))\} \big\|)\ Ai \leq d$
 (is $\bigwedge Ai. Ai \in A \implies out-degree\ (with-proj\ ?G)\ Ai \leq d$)
 <proof>

lemma *obtain-dependency-graph:*

assumes $A \neq \{\}$
assumes $F\ 'A \subseteq events$
assumes *finite* A
assumes $\bigwedge Ai. Ai \in A \implies$
 $(\exists S . S \subseteq A - \{Ai\} \wedge card\ S \geq card\ A - d - 1 \wedge mutual-indep-events\ (F\ Ai)\ F\ S)$
obtains G **where** $dependency-digraph\ G\ M\ F\ pverts\ G = A \bigwedge Ai. Ai \in A \implies$
 $out-degree\ G\ Ai \leq d$
 <proof>

This is the variation of the symmetric version most commonly in use

theorem *lovasz-local-symmetric:*

fixes $d :: nat$
assumes $A \neq \{\}$
assumes $F\ 'A \subseteq events$
assumes *finite* A
assumes $\bigwedge Ai. Ai \in A \implies (\exists S . S \subseteq A - \{Ai\} \wedge card\ S \geq card\ A - d - 1$
 $\wedge mutual-indep-events\ (F\ Ai)\ F\ S)$

```

assumes  $\bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq p$ 
assumes  $\exp(1) * p * (d + 1) \leq 1$ 
shows  $\text{prob } (\bigcap Ai \in A . (\text{space } M - (F Ai))) > 0$ 
<proof>

lemma lovasz-local-symmetric4-set:
  fixes  $d :: \text{nat}$ 
  assumes  $A \neq \{\}$ 
  assumes  $F ' A \subseteq \text{events}$ 
  assumes finite A
  assumes  $\bigwedge Ai. Ai \in A \implies (\exists S . S \subseteq A - \{Ai\} \wedge \text{card } S \geq \text{card } A - d - 1$ 
 $\wedge \text{mutual-indep-events } (F Ai) F S)$ 
  assumes  $\bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq p$ 
  assumes  $4 * p * d \leq 1$ 
  assumes  $d \geq 1$ 
  shows  $\text{prob } (\bigcap Ai \in A . (\text{space } M - F Ai)) > 0$ 
<proof>
end

end
theory Lovasz-Local-Root
  imports
    PiE-Rel-Extras
    Digraph-Extensions

    Prob-Events-Extras
    Cond-Prob-Extensions
    Indep-Events

    Basic-Method
    Lovasz-Local-Lemma
begin
end

```

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