

Kraus Maps*

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Abstract

We formalize Kraus maps [1, Section 3], i.e., quantum channels of the form $\rho \mapsto \sum_x M_x \rho M_x^\dagger$ for suitable families of “Kraus operators” M_x . (In the finite-dimensional setting and the setting of separable Hilbert spaces, those are known to be equivalent to completely-positive maps, another common formalization of quantum channels.) Our results hold for arbitrary (i.e., not necessarily finite-dimensional or separable) Hilbert spaces.

Specifically, in theory `Kraus_Families`, we formalize the type (α, β, ξ) `kraus_family` of families of Kraus operators M_x (Kraus families for short), from trace-class operators on Hilbert space α to those on β , indexed by x of type ξ . This induces both a Kraus map $\rho \mapsto \sum_x M_x \rho M_x^\dagger$, as well as a quantum measurement with outcomes of type ξ .

We define and study various special Kraus families such as zero, identity, application of an operator, sum of two Kraus maps, sequential composition, infinite sum, random sampling, trace, tensor products, and complete measurements.

Furthermore, since working with explicit Kraus families can be cumbersome when the specific family of operators is not relevant, in theory `Kraus_Maps`, we define a Kraus map to be a function between two spaces that is of the form $\rho \mapsto \sum_x M_x \rho M_x^\dagger$ for some Kraus family, and restate our results in terms of such functions.

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1 Miscellaneous missing theorems

theory *Misc-Kraus-Maps*

imports

Hilbert-Space-Tensor-Product.Hilbert-Space-Tensor-Product

Hilbert-Space-Tensor-Product.Von-Neumann-Algebras

begin

unbundle *cblinfun-syntax*

lemma *abs-summable-norm*:

assumes $\langle f \text{ abs-summable-on } A \rangle$

shows $\langle (\lambda x. \text{norm } (f x)) \text{ abs-summable-on } A \rangle$

$\langle \text{proof} \rangle$

lemma *abs-summable-on-add*:

assumes $\langle f \text{ abs-summable-on } A \rangle$ **and** $\langle g \text{ abs-summable-on } A \rangle$

shows $\langle (\lambda x. f x + g x) \text{ abs-summable-on } A \rangle$

$\langle \text{proof} \rangle$

lemma *bdd-above-transform-mono-pos*:

assumes *bdd*: $\langle \text{bdd-above } ((\lambda x. g x) \text{ ' } M) \rangle$

assumes *gpos*: $\langle \bigwedge x. x \in M \implies g x \geq 0 \rangle$

assumes *mono*: $\langle \text{mono-on } (\text{Collect } ((\leq) 0)) f \rangle$

shows $\langle \text{bdd-above } ((\lambda x. f (g x)) \text{ ' } M) \rangle$

$\langle \text{proof} \rangle$

lemma *Ex-iffI*:

assumes $\langle \bigwedge x. P\ x \implies Q\ (f\ x) \rangle$

assumes $\langle \bigwedge x. Q\ x \implies P\ (g\ x) \rangle$

shows $\langle \text{Ex}\ P \longleftrightarrow \text{Ex}\ Q \rangle$

$\langle \text{proof} \rangle$

lemma *has-sum-Sigma'-banach*:

fixes $A :: 'a\ \text{set}$ **and** $B :: 'a \Rightarrow 'b\ \text{set}$

and $f :: 'a \Rightarrow 'b \Rightarrow 'c::\text{banach}$

assumes $\langle (\lambda(x,y). f\ x\ y)\ \text{has-sum}\ S \rangle\ (Sigma\ A\ B)$

shows $\langle (\lambda x. \text{infsum}\ (f\ x)\ (B\ x))\ \text{has-sum}\ S \rangle\ A$

$\langle \text{proof} \rangle$

lemma *infsum-Sigma-topological-monoid*:

fixes $A :: 'a\ \text{set}$ **and** $B :: 'a \Rightarrow 'b\ \text{set}$

and $f :: 'a \times 'b \Rightarrow 'c::\{\text{topological-comm-monoid-add},\ t3\text{-space}\}$

assumes *summableAB*: $f\ \text{summable-on}\ (Sigma\ A\ B)$

assumes *summableB*: $\langle \bigwedge x. x \in A \implies (\lambda y. f\ (x, y))\ \text{summable-on}\ (B\ x) \rangle$

shows $\text{infsum}\ f\ (Sigma\ A\ B) = (\sum_{x \in A} \sum_{y \in B\ x} f\ (x, y))$

$\langle \text{proof} \rangle$

lemma *flip-eq-const*: $\langle (\lambda y. y = x) = ((=)\ x) \rangle$

$\langle \text{proof} \rangle$

lemma *sgn-ket[simp]*: $\langle \text{sgn}\ (ket\ x) = ket\ x \rangle$

$\langle \text{proof} \rangle$

lemma *tensor-op-in-tensor-vn*:

assumes $\langle a \in A \rangle$ **and** $\langle b \in B \rangle$

shows $\langle a \otimes_o b \in A \otimes_{vN} B \rangle$

$\langle \text{proof} \rangle$

lemma *commutant-tensor-vn-subset*:

assumes $\langle \text{von-neumann-algebra}\ A \rangle$ **and** $\langle \text{von-neumann-algebra}\ B \rangle$

shows $\langle \text{commutant}\ A \otimes_{vN} \text{commutant}\ B \subseteq \text{commutant}\ (A \otimes_{vN} B) \rangle$

$\langle \text{proof} \rangle$

lemma *commutant-span[simp]*: $\langle \text{commutant}\ (\text{span}\ X) = \text{commutant}\ X \rangle$

$\langle \text{proof} \rangle$

lemma *explicit-cblinfun-exists-0[simp]*: $\langle \text{explicit-cblinfun-exists}\ (\lambda - -. 0) \rangle$

$\langle \text{proof} \rangle$

lemma *explicit-cblinfun-0[simp]*: $\langle \text{explicit-cblinfun}\ (\lambda - -. 0) = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *cnj-of-bool[simp]*: $\langle \text{cnj } (\text{of-bool } b) = \text{of-bool } b \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-single*:
fixes $f :: \langle - \Rightarrow - :: \{\text{comm-monoid-add}, \text{t2-space}\} \rangle$
assumes $\bigwedge j. j \neq i \implies j \in A \implies f j = 0$
assumes $\langle s = (\text{if } i \in A \text{ then } f i \text{ else } 0) \rangle$
shows $\text{HAS-SUM } f A s$
 $\langle \text{proof} \rangle$

lemma *classical-operator-None[simp]*: $\langle \text{classical-operator } (\lambda -. \text{None}) = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-in-in-closedsubspace*:
assumes $\langle \text{has-sum-in } T f A l \rangle$
assumes $\langle \bigwedge x. x \in A \implies f x \in S \rangle$
assumes $\langle \text{closedin } T S \rangle$
assumes $\langle \text{csubspace } S \rangle$
shows $\langle l \in S \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-coordinatewise*:
 $\langle (f \text{ has-sum } s) A \longleftrightarrow (\forall i. ((\lambda x. f x i) \text{ has-sum } s i) A) \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-butterfly*:
 $\langle \text{butterfly } g h = (\text{one-dim-iso } g * \text{cnj } (\text{one-dim-iso } h)) *_C 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-tc-butterfly*:
fixes $g :: \langle 'a :: \text{one-dim} \rangle$ **and** $h :: \langle 'b :: \text{one-dim} \rangle$
shows $\langle \text{tc-butterfly } g h = (\text{one-dim-iso } g * \text{cnj } (\text{one-dim-iso } h)) *_C 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-iso-of-real[simp]*: $\langle \text{one-dim-iso } (\text{of-real } x) = \text{of-real } x \rangle$
 $\langle \text{proof} \rangle$

lemma *filter-insert-if*:
 $\langle \text{Set.filter } P (\text{insert } x M) = (\text{if } P x \text{ then } \text{insert } x (\text{Set.filter } P M) \text{ else } \text{Set.filter } P M) \rangle$
 $\langle \text{proof} \rangle$

lemma *filter-empty[simp]*: $\langle \text{Set.filter } P \{\} = \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-in-cong-neutral*:

fixes $f\ g :: \langle 'a \Rightarrow 'b :: \text{comm-monoid-add} \rangle$
assumes $\langle \bigwedge x. x \in T - S \implies g\ x = 0 \rangle$
assumes $\langle \bigwedge x. x \in S - T \implies f\ x = 0 \rangle$
assumes $\langle \bigwedge x. x \in S \cap T \implies f\ x = g\ x \rangle$
shows $\text{has-sum-in } X\ f\ S\ x \longleftrightarrow \text{has-sum-in } X\ g\ T\ x$
 $\langle \text{proof} \rangle$

lemma *infsum-in-cong-neutral*:
fixes $f\ g :: \langle 'a \Rightarrow 'b :: \text{comm-monoid-add} \rangle$
assumes $\langle \bigwedge x. x \in T - S \implies g\ x = 0 \rangle$
assumes $\langle \bigwedge x. x \in S - T \implies f\ x = 0 \rangle$
assumes $\langle \bigwedge x. x \in S \cap T \implies f\ x = g\ x \rangle$
shows $\langle \text{infsum-in } X\ f\ S = \text{infsum-in } X\ g\ T \rangle$
 $\langle \text{proof} \rangle$

lemma *filter-image*: $\langle \text{Set.filter } P\ (f\ ' X) = f\ ' (\text{Set.filter } (\lambda x. P\ (f\ x))\ X) \rangle$
 $\langle \text{proof} \rangle$

lemma *Sigma-image-left*: $\langle (\text{SIGMA } x:f\ A. B\ x) = (\lambda(x,y). (f\ x, y))\ ' (\text{SIGMA } x:A. B\ (f\ x)) \rangle$
 $\langle \text{proof} \rangle$

lemma *finite-subset-filter-image*:
assumes *finite* B
assumes $\langle B \subseteq \text{Set.filter } P\ (f\ ' A) \rangle$
shows $\exists C \subseteq A. \text{finite } C \wedge B = f\ ' C$
 $\langle \text{proof} \rangle$

definition $\langle \text{card-le-1 } M \longleftrightarrow (\exists x. M \subseteq \{x\}) \rangle$

lemma *card-le-1-empty[iff]*: $\langle \text{card-le-1 } \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-signleton[iff]*: $\langle \text{card-le-1 } \{x\} \rangle$
 $\langle \text{proof} \rangle$

lemma *sgn-tensor-ell2*: $\langle \text{sgn } (h \otimes_s k) = \text{sgn } h \otimes_s \text{sgn } k \rangle$
 $\langle \text{proof} \rangle$

lemma *is-ortho-set-tensor*:
assumes $\langle \text{is-ortho-set } B \rangle$
assumes $\langle \text{is-ortho-set } C \rangle$
shows $\langle \text{is-ortho-set } ((\lambda(x, y). x \otimes_s y)\ ' (B \times C)) \rangle$
 $\langle \text{proof} \rangle$

end

2 Kraus families

theory *Kraus-Families*

imports

Wlog.Wlog

Hilbert-Space-Tensor-Product.Partial-Trace

Misc-Kraus-Maps

abbrevs

$=_{kr} = =_{kr}$ **and** $=_{kr} = \equiv_{kr}$ **and** $*_{kr} = *__{kr}$

begin

unbundle *cblinfun-syntax*

2.1 Kraus families

definition $\langle \textit{kraus-family } \mathfrak{E} \longleftrightarrow \textit{bdd-above } ((\lambda F. \sum (E,x) \in F. E * o_{CL} E) \text{ ' } \{F. \textit{finite } F \wedge F \subseteq \mathfrak{E}\}) \wedge 0 \notin \textit{fst ' } \mathfrak{E} \rangle$

for $\mathfrak{E} :: \langle ((\textit{--chilbert-space} \Rightarrow_{CL} \textit{--chilbert-space}) \times \textit{--}) \textit{ set} \rangle$

typedef (**overloaded**) $(\textit{'a} :: \textit{chilbert-space}, \textit{'b} :: \textit{chilbert-space}, \textit{'x}) \textit{kraus-family} =$

$\langle \textit{Collect kraus-family} :: ((\textit{'a} \Rightarrow_{CL} \textit{'b}) \times \textit{'x}) \textit{ set set} \rangle$

$\langle \textit{proof} \rangle$

setup-lifting *type-definition-kraus-family*

lemma *kraus-familyI*:

assumes $\langle \textit{bdd-above } ((\lambda F. \sum (E,x) \in F. E * o_{CL} E) \text{ ' } \{F. \textit{finite } F \wedge F \subseteq \mathfrak{E}\}) \rangle$

assumes $\langle 0 \notin \textit{fst ' } \mathfrak{E} \rangle$

shows $\langle \textit{kraus-family } \mathfrak{E} \rangle$

$\langle \textit{proof} \rangle$

lift-definition *kf-apply* $:: \langle (\textit{'a} :: \textit{chilbert-space}, \textit{'b} :: \textit{chilbert-space}, \textit{'x}) \textit{kraus-family} \Rightarrow (\textit{'a}, \textit{'a}) \textit{trace-class} \Rightarrow (\textit{'b}, \textit{'b}) \textit{trace-class} \rangle$ **is**

$\langle \lambda \mathfrak{E} \varrho. (\sum_{\infty} E \in \mathfrak{E}. \textit{sandwich-tc } (\textit{fst } E) \varrho) \rangle$ $\langle \textit{proof} \rangle$

notation *kf-apply* (**infixr** $\langle *__{kr} \rangle$ 70)

lemma *kraus-family-if-finite*[*iff*]: $\langle \textit{kraus-family } \mathfrak{E} \rangle$ **if** $\langle \textit{finite } \mathfrak{E} \rangle$ **and** $\langle 0 \notin \textit{fst ' } \mathfrak{E} \rangle$

$\langle \textit{proof} \rangle$

lemma *kf-apply-scaleC*:

shows $\langle \textit{kf-apply } \mathfrak{E} (c *_C x) = c *_C \textit{kf-apply } \mathfrak{E} x \rangle$

$\langle \textit{proof} \rangle$

lemma *kf-apply-abs-summable*:

assumes $\langle \textit{kraus-family } \mathfrak{E} \rangle$

shows $\langle (\lambda (E,x). \textit{sandwich-tc } E \varrho) \textit{ abs-summable-on } \mathfrak{E} \rangle$

$\langle \textit{proof} \rangle$

lemma *kf-apply-summable*:

shows $\langle (\lambda(E,x). \text{ sandwich-tc } E \ \varrho) \text{ summable-on } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-has-sum*:

shows $\langle ((\lambda(E,x). \text{ sandwich-tc } E \ \varrho) \text{ has-sum } \text{kf-apply } \mathfrak{E} \ \varrho) (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-plus-right*:

shows $\langle \text{kf-apply } \mathfrak{E} (x + y) = \text{kf-apply } \mathfrak{E} x + \text{kf-apply } \mathfrak{E} y \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-uminus-right*:

shows $\langle \text{kf-apply } \mathfrak{E} (-x) = - \text{kf-apply } \mathfrak{E} x \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-minus-right*:

shows $\langle \text{kf-apply } \mathfrak{E} (x - y) = \text{kf-apply } \mathfrak{E} x - \text{kf-apply } \mathfrak{E} y \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-pos*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-apply } \mathfrak{E} \ \varrho \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-mono-right*:

assumes $\langle \varrho \geq \tau \rangle$
shows $\langle \text{kf-apply } \mathfrak{E} \ \varrho \geq \text{kf-apply } \mathfrak{E} \ \tau \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-geq-sum*:

assumes $\langle \varrho \geq 0 \rangle$ **and** $\langle M \subseteq \text{Rep-kraus-family } \mathfrak{E} \rangle$
shows $\langle \text{kf-apply } \mathfrak{E} \ \varrho \geq (\sum (E,-) \in M. \text{ sandwich-tc } E \ \varrho) \rangle$
 $\langle \text{proof} \rangle$

lift-definition *kf-domain* :: $\langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow 'x \text{ set} \rangle$ **is**
 $\langle \lambda \mathfrak{E}. \text{ snd } ' \mathfrak{E} \rangle \langle \text{proof} \rangle$

lemma *kf-apply-clinear[iff]*: $\langle \text{clinear } (\text{kf-apply } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-0-right[iff]*: $\langle \text{kf-apply } \mathfrak{E} \ 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lift-definition $kf\text{-operators} :: \langle ('a::chilbert\text{-space}, 'b::chilbert\text{-space}, 'x) \text{ kraus-family} \Rightarrow ('a \Rightarrow_{CL} 'b) \text{ set} \rangle$ **is**
 $\langle image\ fst :: ('a \Rightarrow_{CL} 'b \times 'x) \text{ set} \Rightarrow ('a \Rightarrow_{CL} 'b) \text{ set} \rangle \langle proof \rangle$

2.2 Bound and norm

lift-definition $kf\text{-bound} :: \langle ('a::chilbert\text{-space}, 'b::chilbert\text{-space}, 'x) \text{ kraus-family} \Rightarrow ('a \Rightarrow_{CL} 'b) \rangle$ **is**
 $\langle \lambda \mathfrak{E}. \text{infsum-in cweak-operator-topology } (\lambda(E,x). E * o_{CL} E) \ \mathfrak{E} \rangle \langle proof \rangle$

lemma $kf\text{-bound-def}'$:

$\langle kf\text{-bound } \mathfrak{E} = \text{Rep-cblinfun-wot } (\sum_{\infty} (E,x) \in \text{Rep-kraus-family } \mathfrak{E}. \text{compose-wot } (\text{adj-wot } (\text{Abs-cblinfun-wot } E)) (\text{Abs-cblinfun-wot } E)) \rangle \langle proof \rangle$

definition $\langle kf\text{-norm } \mathfrak{E} = \text{norm } (kf\text{-bound } \mathfrak{E}) \rangle$

lemma $kf\text{-norm-sum-bdd}$: $\langle bdd\text{-above } ((\lambda F. \text{norm } (\sum (E,x) \in F. E * o_{CL} E)) \text{ ' } \{F. F \subseteq \text{Rep-kraus-family } \mathfrak{E} \wedge \text{finite } F\}) \rangle \langle proof \rangle$

lemma $kf\text{-norm-geq0}$ [iff]:

shows $\langle kf\text{-norm } \mathfrak{E} \geq 0 \rangle$
 $\langle proof \rangle$

lemma $kf\text{-bound-has-sum}$:

shows $\langle \text{has-sum-in cweak-operator-topology } (\lambda(E,x). E * o_{CL} E) (\text{Rep-kraus-family } \mathfrak{E}) (kf\text{-bound } \mathfrak{E}) \rangle \langle proof \rangle$

lemma $\text{kraus-family-iff-summable}$:

$\langle \text{kraus-family } \mathfrak{E} \longleftrightarrow \text{summable-on-in cweak-operator-topology } (\lambda(E,x). E * o_{CL} E) \ \mathfrak{E} \wedge 0 \notin \text{fst ' } \mathfrak{E} \rangle \langle proof \rangle$

lemma $\text{kraus-family-iff-summable}'$:

$\langle \text{kraus-family } \mathfrak{E} \longleftrightarrow (\lambda(E,x). \text{Abs-cblinfun-wot } (E * o_{CL} E)) \text{ summable-on } \mathfrak{E} \wedge 0 \notin \text{fst ' } \mathfrak{E} \rangle \langle proof \rangle$

lemma $kf\text{-bound-summable}$:

shows $\langle \text{summable-on-in cweak-operator-topology } (\lambda(E,x). E * o_{CL} E) (\text{Rep-kraus-family } \mathfrak{E}) \rangle \langle proof \rangle$

lemma $kf\text{-bound-has-sum}'$:

shows $\langle ((\lambda(E,x). \text{compose-wot } (\text{adj-wot } (\text{Abs-cblinfun-wot } E)) (\text{Abs-cblinfun-wot } E)) \text{ has-sum } \text{Abs-cblinfun-wot } (kf\text{-bound } \mathfrak{E})) (\text{Rep-kraus-family } \mathfrak{E}) \rangle \langle proof \rangle$

lemma *kf-bound-summable'*:

$\langle ((\lambda(E,x). \text{compose-wot } (\text{adj-wot } (\text{Abs-cblinfun-wot } E)) (\text{Abs-cblinfun-wot } E)) \text{ summable-on } \text{Rep-kraus-family } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-is-Sup*:

shows $\langle \text{is-Sup } ((\lambda F. \sum (E,x) \in F. E^* \circ_{CL} E) \text{ ' } \{F. \text{finite } F \wedge F \subseteq \text{Rep-kraus-family } \mathfrak{E}\}) \rangle$
 $\langle \text{kf-bound } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-leqI*:

assumes $\langle \bigwedge F. \text{finite } F \implies F \subseteq \text{Rep-kraus-family } \mathfrak{E} \implies (\sum (E,x) \in F. E^* \circ_{CL} E) \leq B \rangle$
shows $\langle \text{kf-bound } \mathfrak{E} \leq B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-pos[iff]*: $\langle \text{kf-bound } \mathfrak{E} \geq 0 \rangle$

$\langle \text{proof} \rangle$

lemma *not-not-singleton-kf-norm-0*:

fixes $\mathfrak{E} :: \langle ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
assumes $\langle \neg \text{class.not-singleton TYPE('a)} \rangle$
shows $\langle \text{kf-norm } \mathfrak{E} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-sum-leqI*:

assumes $\langle \bigwedge F. \text{finite } F \implies F \subseteq \text{Rep-kraus-family } \mathfrak{E} \implies \text{norm } (\sum (E,x) \in F. E^* \circ_{CL} E) \leq B \rangle$
shows $\langle \text{kf-norm } \mathfrak{E} \leq B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-geq-sum*:

assumes $\langle M \subseteq \text{Rep-kraus-family } \mathfrak{E} \rangle$
shows $\langle (\sum (E,-) \in M. E^* \circ_{CL} E) \leq \text{kf-bound } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-geq-norm-sum*:

assumes $\langle M \subseteq \text{Rep-kraus-family } \mathfrak{E} \rangle$
shows $\langle \text{norm } (\sum (E,-) \in M. E^* \circ_{CL} E) \leq \text{kf-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-finite*: $\langle \text{kf-bound } \mathfrak{E} = (\sum (E,x) \in \text{Rep-kraus-family } \mathfrak{E}. E^* \circ_{CL} E) \rangle$ **if** $\langle \text{finite } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-finite*: $\langle \text{kf-norm } \mathfrak{E} = \text{norm } (\sum (E,x) \in \text{Rep-kraus-family } \mathfrak{E}. E^* \circ_{CL} E) \rangle$

if $\langle \text{finite } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-bounded-pos*:
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{norm } (\text{kf-apply } \mathfrak{E} \ \varrho) \leq \text{kf-norm } \mathfrak{E} * \text{norm } \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-bounded*:
— We suspect the factor 4 is not needed here but don't know how to prove that.
 $\langle \text{norm } (\text{kf-apply } \mathfrak{E} \ \varrho) \leq 4 * \text{kf-norm } \mathfrak{E} * \text{norm } \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-bounded-clinear*[*bounded-clinear*]:
shows $\langle \text{bounded-clinear } (\text{kf-apply } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-from-map*: $\langle \psi \cdot_C \text{kf-bound } \mathfrak{E} \ \varphi = \text{trace-tc } (\mathfrak{E} *_{k_r} \text{tc-butterfly } \varphi \ \psi) \rangle$
 $\langle \text{proof} \rangle$

lemma *trace-from-kf-bound*: $\langle \text{trace-tc } (\mathfrak{E} *_{k_r} \varrho) = \text{trace-tc } (\text{compose-tcr } (\text{kf-bound } \mathfrak{E}) \ \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-selfadjoint*[*iff*]: $\langle \text{selfadjoint } (\text{kf-bound } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-leq-kf-norm-id*:
shows $\langle \text{kf-bound } \mathfrak{E} \leq \text{kf-norm } \mathfrak{E} *_R \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

2.3 Basic Kraus families

lemma *kf-emptyset*[*iff*]: $\langle \text{kraus-family } \{\} \rangle$
 $\langle \text{proof} \rangle$

instantiation *kraus-family* :: (*chilbert-space*, *chilbert-space*, *type*) $\langle \text{zero} \rangle$ **begin**

lift-definition *zero-kraus-family* :: $\langle ('a, 'b, 'x) \text{kraus-family} \rangle$ **is** $\langle \{\} \rangle$

$\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$

end

lemma *kf-apply-0*[*simp*]: $\langle \text{kf-apply } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-0*[*simp*]: $\langle \text{kf-bound } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-0*[*simp*]: $\langle \text{kf-norm } 0 = 0 \rangle$

⟨proof⟩

lift-definition $kf\text{-of-op} :: \langle ('a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \text{chilbert-space}) \Rightarrow ('a, 'b, \text{unit}) \text{ kraus-family} \rangle$
is

⟨ $\lambda E :: 'a \Rightarrow_{CL} 'b. \text{ if } E = 0 \text{ then } \{\} \text{ else } \{(E, ())\}$ ⟩
 ⟨proof⟩

lemma $kf\text{-of-op}0[simp]: \langle kf\text{-of-op } 0 = 0 \rangle$
 ⟨proof⟩

lemma $kf\text{-of-op-norm}[simp]: \langle kf\text{-norm } (kf\text{-of-op } E) = (\text{norm } E)^2 \rangle$
 ⟨proof⟩

lemma $kf\text{-operators-in-kf-of-op}[simp]: \langle kf\text{-operators } (kf\text{-of-op } U) = (\text{if } U = 0 \text{ then } \{\} \text{ else } \{U\}) \rangle$
 ⟨proof⟩

lemma $kf\text{-domain-of-op}[simp]: \langle kf\text{-domain } (kf\text{-of-op } A) = \{()\} \text{ if } \langle A \neq 0 \rangle$
 ⟨proof⟩

definition $\langle kf\text{-id} = kf\text{-of-op id-cblinfun} \rangle$

lemma $kf\text{-domain-id}[simp]: \langle kf\text{-domain } (kf\text{-id} :: ('a :: \{\text{chilbert-space}, \text{not-singleton}\}, -, -) \text{ kraus-family})$
 $= \{()\} \rangle$
 ⟨proof⟩

lemma $kf\text{-of-op-id}[simp]: \langle kf\text{-of-op id-cblinfun} = kf\text{-id} \rangle$
 ⟨proof⟩

lemma $kf\text{-norm-id-leq1}: \langle kf\text{-norm } kf\text{-id} \leq 1 \rangle$
 ⟨proof⟩

lemma $kf\text{-norm-id-eq1}[simp]: \langle kf\text{-norm } (kf\text{-id} :: ('a :: \{\text{chilbert-space}, \text{not-singleton}\}, 'a, \text{unit})$
 $\text{ kraus-family}) = 1 \rangle$
 ⟨proof⟩

lemma $kf\text{-operators-in-kf-id}[simp]: \langle kf\text{-operators } kf\text{-id} = (\text{if } (id\text{-cblinfun} :: 'a :: \text{chilbert-space} \Rightarrow_{CL} -) = 0$
 $\text{ then } \{\} \text{ else } \{id\text{-cblinfun} :: 'a \Rightarrow_{CL} -\}) \rangle$
 ⟨proof⟩

instantiation $\text{ kraus-family} :: (\text{chilbert-space}, \text{chilbert-space}, \text{type}) \text{ scaleR begin}$

lift-definition $\text{ scaleR-kraus-family} :: \langle \text{real} \Rightarrow ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family}$
 $\Rightarrow ('a, 'b, 'x) \text{ kraus-family} \rangle$ **is**

⟨ $\lambda r \in \mathbb{C}. \text{ if } r \leq 0 \text{ then } \{\} \text{ else } (\lambda(E, x). (\text{sqrt } r *_R E, x)) \text{ ' } \mathbb{C} \rangle$

⟨proof⟩

instance ⟨proof⟩

end

lemma $kf\text{-scale-apply}:$

assumes $\langle r \geq 0 \rangle$
shows $\langle \text{kf-apply } (r *_R \mathfrak{E}) \varrho = r *_R \text{kf-apply } \mathfrak{E} \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-scale[simp]*: $\langle \text{kf-bound } (c *_R \mathfrak{E}) = c *_R \text{kf-bound } \mathfrak{E} \rangle$ **if** $\langle c \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-scale[simp]*: $\langle \text{kf-norm } (c *_R \mathfrak{E}) = c * \text{kf-norm } \mathfrak{E} \rangle$ **if** $\langle c \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-of-op-apply*: $\langle \text{kf-apply } (\text{kf-of-op } E) \varrho = \text{sandwich-tc } E \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-id-apply[simp]*: $\langle \text{kf-apply } \text{kf-id } \varrho = \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-scale-apply-neg*:
assumes $\langle r \leq 0 \rangle$
shows $\langle \text{kf-apply } (r *_R \mathfrak{E}) = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-0-left[iff]*: $\langle \text{kf-apply } 0 \varrho = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-of-op[simp]*: $\langle \text{kf-bound } (\text{kf-of-op } A) = A * o_{CL} A \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-id[simp]*: $\langle \text{kf-bound } \text{kf-id} = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

2.4 Filtering

lift-definition *kf-filter* :: $\langle ('x \Rightarrow \text{bool}) \Rightarrow ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda(P::'x \Rightarrow \text{bool}) \mathfrak{E}. \text{Set.filter } (\lambda(E,x). P x) \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-false[simp]*: $\langle \text{kf-filter } (\lambda-. \text{False}) \mathfrak{E} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-true[simp]*: $\langle \text{kf-filter } (\lambda-. \text{True}) \mathfrak{E} = \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

definition $\langle \text{kf-apply-on } \mathfrak{E} S = \text{kf-apply } (\text{kf-filter } (\lambda x. x \in S)) \mathfrak{E} \rangle$

notation *kf-apply-on* $(- *_R @- / - [71, 1000, 70] 70)$

lemma *kf-apply-on-pos*:
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-apply-on } \mathfrak{E} X \varrho \geq 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-bounded-clinear*[*bounded-clinear*]:

shows $\langle \text{bounded-clinear } (kf\text{-apply-on } \mathfrak{E} \ X) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-filter-twice*:

$\langle kf\text{-filter } P \ (kf\text{-filter } Q \ \mathfrak{E}) = kf\text{-filter } (\lambda x. P \ x \wedge Q \ x) \ \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-union-has-sum*:

assumes $\langle \bigwedge X \ Y. X \in F \implies Y \in F \implies X \neq Y \implies \text{disjnt } X \ Y \rangle$

shows $\langle ((\lambda X. kf\text{-apply-on } \mathfrak{E} \ X \ \varrho) \text{ has-sum } (kf\text{-apply-on } \mathfrak{E} \ (\bigcup F) \ \varrho)) \ F \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-empty*[*simp*]: $\langle kf\text{-apply-on } E \ \{\} \ \varrho = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-union-eqI*:

assumes $\langle \bigwedge X \ Y. (X, Y) \in F \implies kf\text{-apply-on } \mathfrak{E} \ X \ \varrho = kf\text{-apply-on } \mathfrak{F} \ Y \ \varrho \rangle$

assumes $\langle \bigwedge X \ Y \ X' \ Y'. (X, Y) \in F \implies (X', Y') \in F \implies (X, Y) \neq (X', Y') \implies \text{disjnt } X \ X' \rangle$

assumes $\langle \bigwedge X \ Y \ X' \ Y'. (X, Y) \in F \implies (X', Y') \in F \implies (X, Y) \neq (X', Y') \implies \text{disjnt } Y \ Y' \rangle$

assumes $XX: \langle XX = \bigcup (fst \ ' F) \rangle$ **and** $YY: \langle YY = \bigcup (snd \ ' F) \rangle$

shows $\langle kf\text{-apply-on } \mathfrak{E} \ XX \ \varrho = kf\text{-apply-on } \mathfrak{F} \ YY \ \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-UNIV*[*simp*]: $\langle kf\text{-apply-on } \mathfrak{E} \ UNIV = kf\text{-apply } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-CARD-1*[*simp*]: $\langle (\lambda \mathfrak{E}. kf\text{-apply-on } \mathfrak{E} \ \{x:::CARD-1\}) = kf\text{-apply} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-union-summable-on*:

assumes $\langle \bigwedge X \ Y. X \in F \implies Y \in F \implies X \neq Y \implies \text{disjnt } X \ Y \rangle$

shows $\langle (\lambda X. kf\text{-apply-on } \mathfrak{E} \ X \ \varrho) \text{ summable-on } F \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-union-infsum*:

assumes $\langle \bigwedge X \ Y. X \in F \implies Y \in F \implies X \neq Y \implies \text{disjnt } X \ Y \rangle$

shows $\langle (\sum_{\infty} X \in F. kf\text{-apply-on } \mathfrak{E} \ X \ \varrho) = kf\text{-apply-on } \mathfrak{E} \ (\bigcup F) \ \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *kf-bound-filter*:

$\langle kf\text{-bound } (kf\text{-filter } P \ \mathfrak{E}) \leq kf\text{-bound } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-filter*:

$\langle \text{kf-norm } (\text{kf-filter } P \ \mathfrak{E}) \leq \text{kf-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-filter[simp]*:

$\langle \text{kf-domain } (\text{kf-filter } P \ E) = \text{Set.filter } P \ (\text{kf-domain } E) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-0-right[simp]*: $\langle \text{kf-filter } P \ 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-0-right[simp]*: $\langle \text{kf-apply-on } \mathfrak{E} \ X \ 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-0-left[simp]*: $\langle \text{kf-apply-on } 0 \ X \ \varrho = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-mono3*:

assumes $\langle \varrho \leq \sigma \rangle$
shows $\langle \mathfrak{E} *_{kr} @X \ \varrho \leq \mathfrak{E} *_{kr} @X \ \sigma \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-on-mono2*:

assumes $\langle X \subseteq Y \rangle$ **and** $\langle \varrho \geq 0 \rangle$
shows $\langle \mathfrak{E} *_{kr} @X \ \varrho \leq \mathfrak{E} *_{kr} @Y \ \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-filter*: $\langle \text{kf-operators } (\text{kf-filter } P \ \mathfrak{E}) \subseteq \text{kf-operators } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-equal-if-filter-equal*:

assumes $\langle \bigwedge x. \text{kf-filter } ((=)x) \ \mathfrak{E} = \text{kf-filter } ((=)x) \ \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} = \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

2.5 Equivalence

definition $\langle \text{kf-eq-weak } \mathfrak{E} \ \mathfrak{F} \longleftrightarrow \text{kf-apply } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \rangle$

definition $\langle \text{kf-eq } \mathfrak{E} \ \mathfrak{F} \longleftrightarrow (\forall x. \text{kf-apply-on } \mathfrak{E} \ \{x\} = \text{kf-apply-on } \mathfrak{F} \ \{x\}) \rangle$

open-bundle *kraus-map-syntax* **begin**

notation *kf-eq-weak* (**infix** $=_{kr}$ 50)

notation *kf-eq* (**infix** $\equiv_{kr}$ 50)

notation *kf-apply* (**infixr** $\langle *_{kr} \rangle$ 70)

notation *kf-apply-on* ($- *_{kr} @-$ / - [71, 1000, 70] 70)

end

lemma *kf-eq-weak-refl*[*iff*]: $\langle x =_{kr} x \rangle$
 $\langle proof \rangle$

lemma *kf-eq-weak-refl0*[*iff*]: $\langle 0 =_{kr} 0 \rangle$
 $\langle proof \rangle$

lemma *kf-bound-cong*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle kf-bound\ \mathfrak{E} = kf-bound\ \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-norm-cong*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle kf-norm\ \mathfrak{E} = kf-norm\ \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-eq-weakI*:
assumes $\langle \bigwedge \varrho. \varrho \geq 0 \implies \mathfrak{E} *_{kr} \varrho = \mathfrak{F} *_{kr} \varrho \rangle$
shows $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-eqI*:
assumes $\langle \bigwedge x. \varrho. \varrho \geq 0 \implies \mathfrak{E} *_{kr} @\{x\} \varrho = \mathfrak{F} *_{kr} @\{x\} \varrho \rangle$
shows $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-eq-weak-trans*[*trans*]:
 $\langle F =_{kr} G \implies G =_{kr} H \implies F =_{kr} H \rangle$
 $\langle proof \rangle$

lemma *kf-apply-eqI*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} *_{kr} \varrho = \mathfrak{F} *_{kr} \varrho \rangle$
 $\langle proof \rangle$

lemma *kf-eq-imp-eq-weak*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-filter-cong-eq*[*cong*]:
assumes $\langle \mathfrak{E} = \mathfrak{F} \rangle$
assumes $\langle \bigwedge x. x \in kf-domain\ \mathfrak{E} \implies P\ x = Q\ x \rangle$
shows $\langle kf-filter\ P\ \mathfrak{E} = kf-filter\ Q\ \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-filter-cong*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{E} \implies P\ x = Q\ x \rangle$
shows $\langle \text{kf-filter } P\ \mathfrak{E} \equiv_{kr} \text{kf-filter } Q\ \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-cong-weak*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{E} \implies P\ x = Q\ x \rangle$
shows $\langle \text{kf-filter } P\ \mathfrak{E} \equiv_{kr} \text{kf-filter } Q\ \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-refl[iff]*: $\langle \mathfrak{E} \equiv_{kr} \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-trans[trans]*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
assumes $\langle \mathfrak{F} \equiv_{kr} \mathfrak{G} \rangle$
shows $\langle \mathfrak{E} \equiv_{kr} \mathfrak{G} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-sym[sym]*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{F} \equiv_{kr} \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-weak-imp-eq-CARD-1*:
fixes $\mathfrak{E}\ \mathfrak{F} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x::\text{CARD-1})\ \text{kraus-family} \rangle$
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-on-eqI-filter*:
assumes $\langle \text{kf-filter } (\lambda x. x \in X)\ \mathfrak{E} \equiv_{kr} \text{kf-filter } (\lambda x. x \in X)\ \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} *_{kr} @X\ \varrho = \mathfrak{F} *_{kr} @X\ \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-on-eqI*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} *_{kr} @X\ \varrho = \mathfrak{F} *_{kr} @X\ \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-eq0I*:
assumes $\langle \mathfrak{E} \equiv_{kr} 0 \rangle$
shows $\langle \mathfrak{E} *_{kr} \varrho = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-weak0-imp-kf-eq0*:

assumes $\langle \mathfrak{E} =_{kr} 0 \rangle$
shows $\langle \mathfrak{E} \equiv_{kr} 0 \rangle$
 $\langle proof \rangle$

lemma *kf-apply-on-eq0I*:
assumes $\langle \mathfrak{E} =_{kr} 0 \rangle$
shows $\langle \mathfrak{E} *_{kr} @X \varrho = 0 \rangle$
 $\langle proof \rangle$

lemma *kf-filter-to-domain[simp]*:
 $\langle kf\text{-filter } (\lambda x. x \in kf\text{-domain } \mathfrak{E}) \ \mathfrak{E} = \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *kf-eq-0-iff-eq-0*: $\langle E =_{kr} 0 \longleftrightarrow E = 0 \rangle$
 $\langle proof \rangle$

lemma *in-kf-domain-iff-apply-nonzero*:
 $\langle x \in kf\text{-domain } \mathfrak{E} \longleftrightarrow kf\text{-apply-on } \mathfrak{E} \{x\} \neq 0 \rangle$
 $\langle proof \rangle$

lemma *kf-domain-cong*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle kf\text{-domain } \mathfrak{E} = kf\text{-domain } \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-eq-weak-sym[sym]*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{F} =_{kr} \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *kf-eqI-from-filter-eq-weak*:
assumes $\langle \bigwedge x. kf\text{-filter } ((=) \ x) \ E =_{kr} kf\text{-filter } ((=) \ x) \ F \rangle$
shows $\langle E \equiv_{kr} F \rangle$
 $\langle proof \rangle$

lemma *kf-eq-weak-from-separatingI*:
fixes $E :: \langle ('q :: \text{chilbert-space}, 'r :: \text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
and $F :: \langle ('q, 'r, 'y) \text{ kraus-family} \rangle$
assumes $\langle separating\text{-set } (bounded\text{-clinear} :: ((('q, 'q) \text{ trace-class} \Rightarrow ('r, 'r) \text{ trace-class}) \Rightarrow bool) \ S \rangle$
assumes $\langle \bigwedge \varrho. \varrho \in S \Longrightarrow E *_{kr} \varrho = F *_{kr} \varrho \rangle$
shows $\langle E =_{kr} F \rangle$
 $\langle proof \rangle$

lemma *kf-eq-weak-eq-trans[trans]*: $\langle a =_{kr} b \Longrightarrow b \equiv_{kr} c \Longrightarrow a =_{kr} c \rangle$

$\langle \text{proof} \rangle$

lemma *kf-eq-eq-weak-trans*[*trans*]: $\langle a \equiv_{kr} b \implies b \equiv_{kr} c \implies a \equiv_{kr} c \rangle$
 $\langle \text{proof} \rangle$

instantiation *kraus-family* :: (*chilbert-space, chilbert-space, type*) *preorder* **begin**

definition *less-eq-kraus-family* **where** $\langle \mathfrak{E} \leq \mathfrak{F} \longleftrightarrow (\forall x. \forall \varrho \geq 0. \text{kf-apply-on } \mathfrak{E} \{x\} \varrho \leq \text{kf-apply-on } \mathfrak{F} \{x\} \varrho) \rangle$

definition *less-kraus-family* **where** $\langle \mathfrak{E} < \mathfrak{F} \longleftrightarrow \mathfrak{E} \leq \mathfrak{F} \wedge \neg \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$

lemma *kf-antisym*: $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \longleftrightarrow \mathfrak{E} \leq \mathfrak{F} \wedge \mathfrak{F} \leq \mathfrak{E} \rangle$

for $\mathfrak{E} \mathfrak{F} :: \langle ('a, 'b, 'c) \text{ kraus-family} \rangle$

$\langle \text{proof} \rangle$

instance

$\langle \text{proof} \rangle$

end

lemma *kf-apply-on-mono1*:

assumes $\langle \mathfrak{E} \leq \mathfrak{F} \rangle$ **and** $\langle \varrho \geq 0 \rangle$

shows $\langle \mathfrak{E} *_{kr} @X \varrho \leq \mathfrak{F} *_{kr} @X \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-mono-left*: $\langle \mathfrak{E} \leq \mathfrak{F} \implies \varrho \geq 0 \implies \mathfrak{E} *_{kr} \varrho \leq \mathfrak{F} *_{kr} \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-mono*:

assumes $\langle \varrho \geq 0 \rangle$

assumes $\langle \mathfrak{E} \leq \mathfrak{F} \rangle$ **and** $\langle \varrho \leq \sigma \rangle$

shows $\langle \mathfrak{E} *_{kr} \varrho \leq \mathfrak{F} *_{kr} \sigma \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-mono*:

assumes $\langle \varrho \geq 0 \rangle$

assumes $\langle \mathfrak{E} \leq \mathfrak{F} \rangle$ **and** $\langle X \subseteq Y \rangle$ **and** $\langle \varrho \leq \sigma \rangle$

shows $\langle \mathfrak{E} *_{kr} @X \varrho \leq \mathfrak{F} *_{kr} @Y \sigma \rangle$

$\langle \text{proof} \rangle$

lemma *kf-one-dim-is-id*[*simp*]:

fixes $\mathfrak{E} :: \langle ('a::\text{one-dim}, 'a::\text{one-dim}, 'x) \text{ kraus-family} \rangle$

shows $\langle \mathfrak{E} \equiv_{kr} \text{kf-norm } \mathfrak{E} *_R \text{kf-id} \rangle$

$\langle \text{proof} \rangle$

2.6 Mapping and flattening

definition *kf-similar-elements* :: $\langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow ('a \Rightarrow_{CL} 'b) \Rightarrow ('a \Rightarrow_{CL} 'b \times 'x) \text{ set} \rangle$ **where**

— All elements of the Kraus family that are equal up to rescaling (and belong to the same output)

$\langle \text{kf-similar-elements } \mathfrak{E} \ E = \{(F, x) \in \text{Rep-kraus-family } \mathfrak{E}. (\exists r > 0. E = r *_R F)\} \rangle$

definition *kf-element-weight* **where**

— The total weight (norm of the square) of all similar elements

$\langle \text{kf-element-weight } \mathfrak{E} \ E = (\sum_{\infty} (F, -) \in \text{kf-similar-elements } \mathfrak{E} \ E. \text{ norm } (F * o_{CL} F)) \rangle$

lemma *kf-element-weight-geq0[simp]*: $\langle \text{kf-element-weight } \mathfrak{E} \ E \geq 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-similar-elements-abs-summable*:

fixes $\mathfrak{E} :: \langle ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} \rangle$

shows $\langle (\lambda(F, -). F * o_{CL} F) \text{ abs-summable-on } (\text{kf-similar-elements } \mathfrak{E} \ E) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-similar-elements-kf-operators*:

assumes $\langle (F, x) \in \text{kf-similar-elements } \mathfrak{E} \ E \rangle$

shows $\langle F \in \text{span } (\text{kf-operators } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-neq0*: $\langle \text{kf-element-weight } \mathfrak{E} \ E \neq 0 \rangle$

if $\langle (E, x) \in \text{Rep-kraus-family } \mathfrak{E} \rangle$ **and** $\langle E \neq 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-0-left[simp]*: $\langle \text{kf-element-weight } 0 \ E = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-0-right[simp]*: $\langle \text{kf-element-weight } E \ 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-scale*:

assumes $\langle r > 0 \rangle$

shows $\langle \text{kf-element-weight } \mathfrak{E} \ (r *_R E) = \text{kf-element-weight } \mathfrak{E} \ E \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-kf-operators*:

assumes $\langle \text{kf-element-weight } \mathfrak{E} \ E \neq 0 \rangle$

shows $\langle E \in \text{span } (\text{kf-operators } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-map-aux*:

fixes $f :: \langle 'x \Rightarrow 'y \rangle$ **and** $\mathfrak{E} :: \langle ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} \rangle$

defines $\langle B \equiv \text{kf-bound } \mathfrak{E} \rangle$

defines $\langle \text{filtered } y \equiv \text{kf-filter } (\lambda x. f \ x = y) \ \mathfrak{E} \rangle$

defines $\langle \text{flattened} \equiv \{ (E, y). \text{ norm } (E * o_{CL} E) = \text{kf-element-weight } (\text{filtered } y) \ E \wedge E \neq 0 \} \rangle$

defines $\langle \text{good} \equiv (\lambda(E, y). (\text{norm } E)^2 = \text{kf-element-weight } (\text{filtered } y) \ E \wedge E \neq 0) \rangle$

shows $\langle \text{has-sum-in cweak-operator-topology } (\lambda(E, -). E * o_{CL} E) \text{ flattened } B \rangle$ **(is ?has-sum)**

and $\langle \text{snd } ' (\text{SIGMA } (E, y) : \text{Collect good. kf-similar-elements } (\text{filtered } y) \ E) \rangle$

$= \{ (F, x). (F, x) \in \text{Rep-kraus-family } \mathfrak{E} \wedge F \neq 0 \} \rangle$ **(is ?snd-sigma)**

and $\langle \text{inj-on snd } (\text{SIGMA } p : \text{Collect good. kf-similar-elements } (\text{filtered } (\text{snd } p)) \ (\text{fst } p)) \rangle$ **(is**

?inj-snd)
 ⟨proof⟩

lift-definition $\text{kf-map} :: \langle ('x \Rightarrow 'y) \Rightarrow ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow ('a, 'b, 'y) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda f \ \mathfrak{E}. \{(E, y). \text{norm } (E * o_{CL} E) = \text{kf-element-weight } (\text{kf-filter } (\lambda x. f x = y) \ \mathfrak{E}) \ E \wedge E \neq 0\} \rangle$
 ⟨proof⟩

lemma
fixes $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
shows $\text{kf-apply-map[simp]}: \langle \text{kf-apply } (\text{kf-map } f \ \mathfrak{E}) = \text{kf-apply } \mathfrak{E} \rangle$
and $\text{kf-map-bound}: \langle \text{kf-bound } (\text{kf-map } f \ \mathfrak{E}) = \text{kf-bound } \mathfrak{E} \rangle$
and $\text{kf-map-norm[simp]}: \langle \text{kf-norm } (\text{kf-map } f \ \mathfrak{E}) = \text{kf-norm } \mathfrak{E} \rangle$
 ⟨proof⟩

abbreviation $\langle \text{kf-flatten} \equiv \text{kf-map } (\lambda -. ()) \rangle$

Like *kf-map*, but with a much simpler definition. However, only makes sense for injective functions.

lift-definition $\text{kf-map-inj} :: \langle ('x \Rightarrow 'y) \Rightarrow ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow ('a, 'b, 'y) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda f \ \mathfrak{E}. (\lambda(E, x). (E, f x)) \ ' \mathfrak{E} \rangle$
 ⟨proof⟩

lemma *kf-element-weight-map-inj*:
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-element-weight } (\text{kf-map-inj } f \ \mathfrak{E}) \ E = \text{kf-element-weight } \mathfrak{E} \ E \rangle$
 ⟨proof⟩

lemma *kf-eq-weak-kf-map-left*: $\langle \text{kf-map } f \ F =_{kr} G \rangle$ **if** $\langle F =_{kr} G \rangle$
 ⟨proof⟩

lemma *kf-eq-weak-kf-map-right*: $\langle F =_{kr} \text{kf-map } f \ G \rangle$ **if** $\langle F =_{kr} G \rangle$
 ⟨proof⟩

lemma *kf-filter-map*:
fixes $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
shows $\langle \text{kf-filter } P \ (\text{kf-map } f \ \mathfrak{E}) = \text{kf-map } f \ (\text{kf-filter } (\lambda x. P \ (f x)) \ \mathfrak{E}) \rangle$
 ⟨proof⟩

lemma *kf-filter-map-inj*:
fixes $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
shows $\langle \text{kf-filter } P \ (\text{kf-map-inj } f \ \mathfrak{E}) = \text{kf-map-inj } f \ (\text{kf-filter } (\lambda x. P \ (f x)) \ \mathfrak{E}) \rangle$
 ⟨proof⟩

lemma *kf-map-kf-map-inj-comp*:
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$

shows $\langle \text{kf-map } g \ (\text{kf-map-inj } f \ \mathfrak{E}) = \text{kf-map } (g \circ f) \ \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-element-weight-eqweak0*:
assumes $\langle \mathfrak{E} =_{kr} 0 \rangle$
shows $\langle \text{kf-element-weight } \mathfrak{E} \ E = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-kf-map-comp*:
assumes $\langle \text{inj-on } g \ (f \text{ ' } \text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-map-inj } g \ (\text{kf-map } f \ \mathfrak{E}) = \text{kf-map } (g \circ f) \ \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-map-inj[simp]*:
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-map-inj } f \ \mathfrak{E} *_{kr} \varrho = \mathfrak{E} *_{kr} \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-eq-kf-map*:
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-map-inj } f \ \mathfrak{E} \equiv_{kr} \text{kf-map } f \ \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-id[simp]*: $\langle \text{kf-map-inj id } \mathfrak{E} = \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-id*: $\langle \text{kf-map id } \mathfrak{E} \equiv_{kr} \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-bound[simp]*:
fixes $\mathfrak{E} :: \langle ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-bound } (\text{kf-map-inj } f \ \mathfrak{E}) = \text{kf-bound } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-norm[simp]*:
fixes $\mathfrak{E} :: \langle ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-norm } (\text{kf-map-inj } f \ \mathfrak{E}) = \text{kf-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-cong-weak*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \text{kf-map } f \ \mathfrak{E} =_{kr} \text{kf-map } g \ \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-flatten-cong-weak*:

assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \text{kf-flatten } \mathfrak{E} =_{kr} \text{kf-flatten } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-flatten-cong*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \text{kf-flatten } \mathfrak{E} \equiv_{kr} \text{kf-flatten } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-twice*:
 $\langle \text{kf-map } f (\text{kf-map } g \ \mathfrak{E}) \equiv_{kr} \text{kf-map } (f \circ g) \ \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-cong*:
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{E} \implies f \ x = g \ x \rangle$
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle \text{kf-map } f \ \mathfrak{E} \equiv_{kr} \text{kf-map } g \ \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-cong-eq*:
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{E} \implies f \ x = g \ x \rangle$
assumes $\langle \mathfrak{E} = \mathfrak{F} \rangle$
shows $\langle \text{kf-map-inj } f \ \mathfrak{E} = \text{kf-map-inj } g \ \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-map[simp]*:
 $\langle \text{kf-domain } (\text{kf-map } f \ \mathfrak{E}) = f \ ` \ \text{kf-domain } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-on-map[simp]*:
 $\langle (\text{kf-map } f \ E) \ *_{kr} \ @X \ \varrho = E \ *_{kr} \ @(f \ - \ ` \ X) \ \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-on-map-inj[simp]*:
assumes $\langle \text{inj-on } f \ ((f \ - \ ` \ X) \cap \text{kf-domain } E) \rangle$
shows $\langle \text{kf-map-inj } f \ E \ *_{kr} \ @X \ \varrho = E \ *_{kr} \ @(f \ - \ ` \ X) \ \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map0[simp]*: $\langle \text{kf-map } f \ 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-kr-eq-weak*:
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-map-inj } f \ \mathfrak{E} =_{kr} \mathfrak{E} \rangle$

⟨proof⟩

lemma *kf-map-inj-0[simp]*: ⟨*kf-map-inj* *f* 0 = 0⟩
 ⟨proof⟩

lemma *kf-domain-map-inj[simp]*: ⟨*kf-domain* (*kf-map-inj* *f* $\mathfrak{E}$) = *f* ‘ *kf-domain* $\mathfrak{E}$ ⟩
 ⟨proof⟩

lemma *kf-operators-kf-map*:
 ⟨*kf-operators* (*kf-map* *f* $\mathfrak{E}$) $\subseteq$ *span* (*kf-operators* $\mathfrak{E}$)⟩
 ⟨proof⟩

lemma *kf-operators-kf-map-inj[simp]*: ⟨*kf-operators* (*kf-map-inj* *f* $\mathfrak{E}$) = *kf-operators* $\mathfrak{E}$ ⟩
 ⟨proof⟩

2.7 Addition

lift-definition *kf-plus* :: ⟨('a::chilbert-space, 'b::chilbert-space, 'x) *kraus-family* $\Rightarrow$ ('a,'b,'y) *kraus-family* $\Rightarrow$ ('a,'b,'x+'y) *kraus-family*⟩ **is**
 ⟨ $\lambda \mathfrak{E} \mathfrak{F}. (\lambda (E, x). (E, \text{Inl } x)) \text{ ‘ } \mathfrak{E} \cup (\lambda (F, y). (F, \text{Inr } y)) \text{ ‘ } \mathfrak{F}$ ⟩
 ⟨proof⟩

instantiation *kraus-family* :: (*chilbert-space*, *chilbert-space*, *type*) **plus** **begin**

definition *plus-kraus-family* **where** ⟨ $\mathfrak{E} + \mathfrak{F} = \text{kf-map } (\lambda xy. \text{case } xy \text{ of } \text{Inl } x \Rightarrow x \mid \text{Inr } y \Rightarrow y)$ (*kf-plus* $\mathfrak{E}$ $\mathfrak{F}$)⟩

instance⟨proof⟩

end

lemma *kf-plus-apply*:
fixes $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('a, 'b, 'y) \text{ kraus-family} \rangle$
shows ⟨*kf-apply* (*kf-plus* $\mathfrak{E}$ $\mathfrak{F}$) $\varrho = \text{kf-apply } \mathfrak{E} \varrho + \text{kf-apply } \mathfrak{F} \varrho$ ⟩
 ⟨proof⟩

lemma *kf-plus-apply'*: ⟨ $(\mathfrak{E} + \mathfrak{F}) *_{kr} \varrho = \mathfrak{E} *_{kr} \varrho + \mathfrak{F} *_{kr} \varrho$ ⟩
 ⟨proof⟩

lemma *kf-plus-0-left[simp]*: ⟨*kf-plus* 0 $\mathfrak{E} = \text{kf-map-inj Inr } \mathfrak{E}$ ⟩
 ⟨proof⟩

lemma *kf-plus-0-right[simp]*: ⟨*kf-plus* $\mathfrak{E}$ 0 = *kf-map-inj Inl* $\mathfrak{E}$ ⟩
 ⟨proof⟩

lemma *kf-plus-0-left'[simp]*: ⟨0 + $\mathfrak{E} \equiv_{kr} \mathfrak{E}$ ⟩
 ⟨proof⟩

lemma *kf-plus-0-right'*: ⟨ $\mathfrak{E} + 0 \equiv_{kr} \mathfrak{E}$ ⟩
 ⟨proof⟩

lemma *kf-plus-bound*: $\langle \text{kf-bound } (\text{kf-plus } \mathfrak{E} \ \mathfrak{F}) = \text{kf-bound } \mathfrak{E} + \text{kf-bound } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-plus-bound'*: $\langle \text{kf-bound } (\mathfrak{E} + \mathfrak{F}) = \text{kf-bound } \mathfrak{E} + \text{kf-bound } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-triangle*: $\langle \text{kf-norm } (\text{kf-plus } \mathfrak{E} \ \mathfrak{F}) \leq \text{kf-norm } \mathfrak{E} + \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-triangle'*: $\langle \text{kf-norm } (\mathfrak{E} + \mathfrak{F}) \leq \text{kf-norm } \mathfrak{E} + \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-plus-map-both*:
 $\langle \text{kf-plus } (\text{kf-map } f \ \mathfrak{E}) (\text{kf-map } g \ \mathfrak{F}) = \text{kf-map } (\text{map-sum } f \ g) (\text{kf-plus } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

2.8 Composition

lemma *kf-comp-dependent-raw-norm-aux*:
fixes $\mathfrak{E} :: \langle 'a \Rightarrow ('e :: \text{chilbert-space}, 'f :: \text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('b :: \text{chilbert-space}, 'e, 'a) \text{ kraus-family} \rangle$
assumes $B: \langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \text{kf-norm } (\mathfrak{E} \ x) \leq B \rangle$
assumes $[\text{simp}]: \langle B \geq 0 \rangle$
assumes $\langle \text{finite } C \rangle$
assumes $C\text{-subset}: \langle C \subseteq (\lambda((F,y), (E,x)). (E \circ_{CL} F, (F,E,y,x))) \text{ ' } (\text{SIGMA } (F,y): \text{Rep-kraus-family } \mathfrak{F}. \text{Rep-kraus-family } (\mathfrak{E} \ y)) \rangle$
shows $\langle (\sum_{(E,x) \in C} (E * \circ_{CL} E) \leq (B * \text{kf-norm } \mathfrak{F}) *_R \text{id-cblinfun}) \rangle$
 $\langle \text{proof} \rangle$

lift-definition *kf-comp-dependent-raw* :: $\langle ('x \Rightarrow ('b :: \text{chilbert-space}, 'c :: \text{chilbert-space}, 'y) \text{ kraus-family}) \Rightarrow ('a :: \text{chilbert-space}, 'b, 'x) \text{ kraus-family} \rangle$
 $\Rightarrow ('a, 'c, ('a \Rightarrow_{CL} 'b) \times ('b \Rightarrow_{CL} 'c) \times 'x \times 'y) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda \mathfrak{E} \ \mathfrak{F}. \text{if bdd-above } ((\text{kf-norm } \circ \ \mathfrak{E}) \text{ ' } \text{kf-domain } \mathfrak{F}) \text{ then} \langle \text{Set.filter } (\lambda(EF, -). EF \neq 0) ((\lambda((F,y), (E :: 'b \Rightarrow_{CL} 'c, x :: 'y)). (E \circ_{CL} F, (F,E,y,x))) \text{ ' } (\text{SIGMA } (F :: 'a \Rightarrow_{CL} 'b, y :: 'x): \text{Rep-kraus-family } \mathfrak{F}. (\text{Rep-kraus-family } (\mathfrak{E} \ y)))) \text{ else } \{\} \rangle \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-norm-leq*:
fixes $\mathfrak{E} :: \langle 'a \Rightarrow ('b :: \text{chilbert-space}, 'c :: \text{chilbert-space}, 'd) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('e :: \text{chilbert-space}, 'b, 'a) \text{ kraus-family} \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \text{kf-norm } (\mathfrak{E} \ x) \leq B \rangle$
assumes $\langle B \geq 0 \rangle$
shows $\langle \text{kf-norm } (\text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F}) \leq B * \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

hide-fact *kf-comp-dependent-raw-norm-aux*

definition $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} = \text{kf-map } (\lambda(F,E,y,x). (y,x)) (\text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F}) \rangle$

definition $\langle \text{kf-comp } \mathfrak{E} \ \mathfrak{F} = \text{kf-comp-dependent } (\lambda\cdot. \mathfrak{E}) \ \mathfrak{F} \rangle$

lemma *kf-comp-dependent-norm-leq*:

assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \text{kf-norm } (\mathfrak{E} \ x) \leq B \rangle$
assumes $\langle B \geq 0 \rangle$
shows $\langle \text{kf-norm } (\text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F}) \leq B * \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-norm-leq*:

shows $\langle \text{kf-norm } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) \leq \text{kf-norm } \mathfrak{E} * \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-apply*:

fixes $\mathfrak{E} :: \langle 'y \Rightarrow ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('c::\text{chilbert-space}, 'a, 'y) \text{ kraus-family} \rangle$
assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \ \text{`kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F} *_{kr} \varrho$
 $= (\sum_{\infty (F,y) \in \text{Rep-kraus-family } \mathfrak{F}. \ \mathfrak{E} \ y *_{kr} \text{ sandwich-tc } F \ \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-apply*:

fixes $\mathfrak{E} :: \langle 'y \Rightarrow ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('c::\text{chilbert-space}, 'a, 'y) \text{ kraus-family} \rangle$
assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \ \text{`kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} *_{kr} \varrho$
 $= (\sum_{\infty (F,y) \in \text{Rep-kraus-family } \mathfrak{F}. \ \mathfrak{E} \ y *_{kr} \text{ sandwich-tc } F \ \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-apply*:

shows $\langle \text{kf-apply } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) = \text{kf-apply } \mathfrak{E} \circ \text{kf-apply } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-cong-weak*: $\langle \text{kf-comp } F \ G =_{kr} \text{kf-comp } F' \ G' \rangle$

if $\langle F =_{kr} F' \rangle$ **and** $\langle G =_{kr} G' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-assoc*:

fixes $\mathfrak{E} :: \langle 'f \Rightarrow ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('g \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
defines $\langle \text{reorder} :: 'a \Rightarrow_{CL} 'c \times 'c \Rightarrow_{CL} 'd \times ('a \Rightarrow_{CL} 'b \times 'b \Rightarrow_{CL} 'c \times 'g \times 'f) \times 'e$
 $\Rightarrow 'a \Rightarrow_{CL} 'b \times 'b \Rightarrow_{CL} 'd \times 'g \times 'b \Rightarrow_{CL} 'c \times 'c \Rightarrow_{CL} 'd \times 'f \times 'e \equiv$
 $\lambda(FG::'a \Rightarrow_{CL} 'c, E::'c \Rightarrow_{CL} 'd, (G::'a \Rightarrow_{CL} 'b, F::'b \Rightarrow_{CL} 'c, g::'g, f::'f), e::'e).$
 $(G, E \ o_{CL} F, g, F, E, f, e) \rangle$
assumes $\langle \text{bdd-above } (\text{range } (\text{kf-norm } o \ \mathfrak{E})) \rangle$
assumes $\langle \text{bdd-above } (\text{range } (\text{kf-norm } o \ \mathfrak{F})) \rangle$
shows $\langle \text{kf-comp-dependent-raw } (\lambda g::'g. \text{kf-comp-dependent-raw } \mathfrak{E} \ (\mathfrak{F} \ g)) \ \mathfrak{G}$
 $= \text{kf-map-inj reorder } (\text{kf-comp-dependent-raw } (\lambda(-,-,f). \ \mathfrak{E} \ f) \ (\text{kf-comp-dependent-raw } \mathfrak{F}$

$\mathfrak{G})\rangle$
 (is $\langle ?lhs = ?rhs \rangle$)
 $\langle proof \rangle$

lemma *kf-filter-comp-dependent*:

fixes $\mathfrak{F} :: \langle 'e \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
 and $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
 assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{F}) \text{ 'kf-domain } \mathfrak{E}) \rangle$
 shows $\langle \text{kf-filter } (\lambda(e,f). F \ e \ f \wedge E \ e) \ (\text{kf-comp-dependent } \mathfrak{F} \ \mathfrak{E})$
 $= \text{kf-comp-dependent } (\lambda e. \text{kf-filter } (F \ e) \ (\mathfrak{F} \ e)) \ (\text{kf-filter } E \ \mathfrak{E}) \rangle$
 $\langle proof \rangle$

lemma *kf-comp-assoc-weak*:

fixes $\mathfrak{E} :: \langle ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
 and $\mathfrak{F} :: \langle ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
 and $\mathfrak{G} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
 shows $\langle \text{kf-comp } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) \ \mathfrak{G} =_{k\tau} \text{kf-comp } \mathfrak{E} \ (\text{kf-comp } \mathfrak{F} \ \mathfrak{G}) \rangle$
 $\langle proof \rangle$

lemma *kf-comp-dependent-raw-cong-left*:

assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
 assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}') \text{ 'kf-domain } \mathfrak{F}) \rangle$
 assumes $\langle \bigwedge x. x \in \text{snd } \text{'Rep-kraus-family } \mathfrak{F} \implies \mathfrak{E} \ x = \mathfrak{E}' \ x \rangle$
 shows $\langle \text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F} = \text{kf-comp-dependent-raw } \mathfrak{E}' \ \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-comp-dependent-cong-left*:

assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
 assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}') \text{ 'kf-domain } \mathfrak{F}) \rangle$
 assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \mathfrak{E} \ x = \mathfrak{E}' \ x \rangle$
 shows $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} = \text{kf-comp-dependent } \mathfrak{E}' \ \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-domain-comp-dependent-raw-subset*:

$\langle \text{kf-domain } (\text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F}) \subseteq \text{UNIV} \times \text{UNIV} \times (\text{SIGMA } x:\text{kf-domain } \mathfrak{F}. \text{kf-domain } (\mathfrak{E} \ x)) \rangle$
 $\langle proof \rangle$

lemma *kf-domain-comp-dependent-subset*:

$\langle \text{kf-domain } (\text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F}) \subseteq (\text{SIGMA } x:\text{kf-domain } \mathfrak{F}. \text{kf-domain } (\mathfrak{E} \ x)) \rangle$
 $\langle proof \rangle$

lemma *kf-domain-comp-subset*: $\langle \text{kf-domain } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) \subseteq \text{kf-domain } \mathfrak{F} \times \text{kf-domain } \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *kf-apply-comp-dependent-cong*:

fixes $\mathfrak{E} :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e1) \text{ kraus-family} \rangle$

and $\mathfrak{E}' :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e2) \text{ kraus-family} \rangle$
and $\mathfrak{F} \mathfrak{F}' :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
assumes $\text{bdd} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\text{bdd}' :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}') \text{ 'kf-domain } \mathfrak{F}') \rangle$
assumes $\langle f \in \text{kf-domain } \mathfrak{F} \implies \text{kf-apply-on } (\mathfrak{E} \ f) \ E = \text{kf-apply-on } (\mathfrak{E}' \ f) \ E' \rangle$
assumes $\langle \text{kf-apply-on } \mathfrak{F} \ \{f\} = \text{kf-apply-on } \mathfrak{F}' \ \{f\} \rangle$
shows $\langle \text{kf-apply-on } (\text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F}) \ (\{f\} \times E) = \text{kf-apply-on } (\text{kf-comp-dependent } \mathfrak{E}' \ \mathfrak{F}') \ (\{f\} \times E') \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-cong-weak:*

fixes $\mathfrak{E} :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e1) \text{ kraus-family} \rangle$
and $\mathfrak{E}' :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e2) \text{ kraus-family} \rangle$
and $\mathfrak{F} \mathfrak{F}' :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
assumes $\text{bdd} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\text{eq} :: \langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \mathfrak{E} \ x =_{kr} \mathfrak{E}' \ x \rangle$
assumes $\langle \mathfrak{F} \equiv_{kr} \mathfrak{F}' \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} =_{kr} \text{kf-comp-dependent } \mathfrak{E}' \ \mathfrak{F}' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-assoc:*

fixes $\mathfrak{E} :: \langle 'g \Rightarrow 'f \Rightarrow ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle 'g \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
assumes $\text{bdd-E} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \text{case-prod } \mathfrak{E}) \text{ ' (SIGMA } x:\text{kf-domain } \mathfrak{G}. \text{kf-domain } (\mathfrak{F} \ x))) \rangle$
assumes $\text{bdd-F} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{F}) \text{ 'kf-domain } \mathfrak{G}) \rangle$
shows $\langle (\text{kf-comp-dependent } (\lambda g. \text{kf-comp-dependent } (\mathfrak{E} \ g) (\mathfrak{F} \ g)) \ \mathfrak{G}) \equiv_{kr} \text{kf-map } (\lambda ((g,f), e). (g,f,e)) (\text{kf-comp-dependent } (\lambda(g,f). \mathfrak{E} \ g \ f) (\text{kf-comp-dependent } \mathfrak{F} \ \mathfrak{G})) \rangle$
 $\langle \text{is } \langle ?lhs \equiv_{kr} ?rhs \rangle \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-assoc-weak:*

fixes $\mathfrak{E} :: \langle 'g \Rightarrow 'f \Rightarrow ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle 'g \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
assumes $\text{bdd-E} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \text{case-prod } \mathfrak{E}) \text{ ' (SIGMA } x:\text{kf-domain } \mathfrak{G}. \text{kf-domain } (\mathfrak{F} \ x))) \rangle$
assumes $\text{bdd-F} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{F}) \text{ 'kf-domain } \mathfrak{G}) \rangle$
shows $\langle \text{kf-comp-dependent } (\lambda g. \text{kf-comp-dependent } (\mathfrak{E} \ g) (\mathfrak{F} \ g)) \ \mathfrak{G} =_{kr} \text{kf-comp-dependent } (\lambda(g,f). \mathfrak{E} \ g \ f) (\text{kf-comp-dependent } \mathfrak{F} \ \mathfrak{G}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-comp-assoc-weak:*

fixes $\mathfrak{E} :: \langle ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle 'g \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$

and $\mathfrak{G} :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'g) \text{ kraus-family} \rangle$
assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{F}) \text{ 'kf-domain } \mathfrak{G}) \rangle$
shows $\langle \text{kf-comp-dependent } (\lambda g. \text{kf-comp } \mathfrak{E} (\mathfrak{F} \ g)) \ \mathfrak{G} =_{kr} \text{kf-comp } \mathfrak{E} (\text{kf-comp-dependent } \mathfrak{F} \ \mathfrak{G}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-comp-dependent-assoc-weak:*

fixes $\mathfrak{E} :: \langle ('f \Rightarrow ('c::\text{hilbert-space}, 'd::\text{hilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'g) \text{ kraus-family} \rangle$
assumes *bdd-E*: $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp } (\text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F}) \ \mathfrak{G} =_{kr} \text{kf-comp-dependent } (\lambda(-,f). \ \mathfrak{E} \ f) (\text{kf-comp } \mathfrak{F} \ \mathfrak{G}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-assoc:*

fixes $\mathfrak{E} :: \langle ('c::\text{hilbert-space}, 'd::\text{hilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'g) \text{ kraus-family} \rangle$
shows $\langle \text{kf-comp } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) \ \mathfrak{G} \equiv_{kr} \text{kf-map } (\lambda((g,f),e). \ (g,f,e)) (\text{kf-comp } \mathfrak{E} (\text{kf-comp } \mathfrak{F} \ \mathfrak{G})) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-cong:*

fixes $\mathfrak{E} \ \mathfrak{E}' :: \langle ('f \Rightarrow ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} \ \mathfrak{F}' :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'f) \text{ kraus-family} \rangle$
assumes *bdd*: $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \mathfrak{E} \ x \equiv_{kr} \mathfrak{E}' \ x \rangle$
assumes $\langle \mathfrak{F} \equiv_{kr} \mathfrak{F}' \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} \equiv_{kr} \text{kf-comp-dependent } \mathfrak{E}' \ \mathfrak{F}' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-cong:*

fixes $\mathfrak{E} \ \mathfrak{E}' :: \langle ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} \ \mathfrak{F}' :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'f) \text{ kraus-family} \rangle$
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{E}' \rangle$
assumes $\langle \mathfrak{F} \equiv_{kr} \mathfrak{F}' \rangle$
shows $\langle \text{kf-comp } \mathfrak{E} \ \mathfrak{F} \equiv_{kr} \text{kf-comp } \mathfrak{E}' \ \mathfrak{F}' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-comp-dependent-raw-of-op:*

shows $\langle \text{kf-bound } (\text{kf-comp-dependent-raw } \mathfrak{E} (\text{kf-of-op } U)) = \text{sandwich } (U*) (\text{kf-bound } (\mathfrak{E} ())) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-comp-dependent-of-op:*

shows $\langle \text{kf-bound } (\text{kf-comp-dependent } \mathfrak{E} (\text{kf-of-op } U)) = \text{sandwich } (U*) (\text{kf-bound } (\mathfrak{E} ())) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-bound-comp-of-op:*

shows $\langle \text{kf-bound } (\text{kf-comp } \mathfrak{E} (\text{kf-of-op } U)) = \text{sandwich } (U*) (\text{kf-bound } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-comp-dependent-of-op-coiso:*

assumes $\langle \text{isometry } (U*) \rangle$

shows $\langle \text{kf-norm } (\text{kf-comp-dependent } \mathfrak{E} (\text{kf-of-op } U)) = \text{kf-norm } (\mathfrak{E} ()) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-comp-of-op-coiso:*

assumes $\langle \text{isometry } (U*) \rangle$

shows $\langle \text{kf-norm } (\text{kf-comp } \mathfrak{E} (\text{kf-of-op } U)) = \text{kf-norm } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-bound-comp-dependend-raw-iso:*

assumes $\langle \text{isometry } U \rangle$

shows $\langle \text{kf-bound } (\text{kf-comp-dependent-raw } (\lambda-. \text{kf-of-op } U) \mathfrak{E}) = \text{kf-bound } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-bound-comp-dependent-iso:*

assumes $\langle \text{isometry } U \rangle$

shows $\langle \text{kf-bound } (\text{kf-comp-dependent } (\lambda-. \text{kf-of-op } U) \mathfrak{E}) = \text{kf-bound } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-bound-comp-iso:*

assumes $\langle \text{isometry } U \rangle$

shows $\langle \text{kf-bound } (\text{kf-comp } (\text{kf-of-op } U) \mathfrak{E}) = \text{kf-bound } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-comp-dependent-iso:*

assumes $\langle \text{isometry } U \rangle$

shows $\langle \text{kf-norm } (\text{kf-comp-dependent } (\lambda-. \text{kf-of-op } U) \mathfrak{E}) = \text{kf-norm } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-comp-iso:*

assumes $\langle \text{isometry } U \rangle$

shows $\langle \text{kf-norm } (\text{kf-comp } (\text{kf-of-op } U) \mathfrak{E}) = \text{kf-norm } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-0-right[simp]:* $\langle \text{kf-comp-dependent-raw } \mathfrak{E} 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-0-left[simp]:* $\langle \text{kf-comp-dependent-raw } 0 \mathfrak{E} = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-comp-dependent-0-left[simp]*: $\langle \text{kf-comp-dependent } (\lambda -. 0) \ E = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-0-right[simp]*: $\langle \text{kf-comp-dependent } E \ 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-0-left[simp]*: $\langle \text{kf-comp } 0 \ E = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-0-right[simp]*: $\langle \text{kf-comp } E \ 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-comp*:
fixes $\mathfrak{F} :: \langle ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \ \text{kraus-family} \rangle$
and $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'e) \ \text{kraus-family} \rangle$
shows $\langle \text{kf-filter } (\lambda(e,f). \ F \ f \wedge \ E \ e) \ (\text{kf-comp } \mathfrak{F} \ \mathfrak{E})$
 $= \text{kf-comp } (\text{kf-filter } F \ \mathfrak{F}) \ (\text{kf-filter } E \ \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-invalid*:
assumes $\langle \neg \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \ ' \ \text{kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-map-left*:
 $\langle \text{kf-comp-dependent } (\lambda x. \ \text{kf-map } (f \ x) \ (E \ x)) \ F$
 $\equiv_{kr} \text{kf-map } (\lambda(x,y). \ (x, f \ x \ y)) \ (\text{kf-comp-dependent } E \ F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-map-right*:
 $\langle \text{kf-comp-dependent } E \ (\text{kf-map } f \ F)$
 $\equiv_{kr} \text{kf-map } (\lambda(x,y). \ (f \ x, y)) \ (\text{kf-comp-dependent } (\lambda x. \ E \ (f \ x)) \ F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-map-inj-right*:
 $\langle \text{kf-comp-dependent-raw } E \ (\text{kf-map-inj } f \ F)$
 $= \text{kf-map-inj } (\lambda(E,F,x,y). \ (E, F, f \ x, y)) \ (\text{kf-comp-dependent-raw } (\lambda x. \ E \ (f \ x)) \ F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-map-inj-right*:
assumes $\langle \text{inj-on } f \ (\text{kf-domain } F) \rangle$
shows $\langle \text{kf-comp-dependent } E \ (\text{kf-map-inj } f \ F)$
 $= \text{kf-map-inj } (\lambda(x,y). \ (f \ x, y)) \ (\text{kf-comp-dependent } (\lambda x. \ E \ (f \ x)) \ F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-map-right-weak*:

$\langle \text{kf-comp-dependent } E \text{ (kf-map } f \text{ } F) \text{ } \equiv_{kr} \text{ kf-comp-dependent } (\lambda x. E \text{ (} f \text{ } x)) \text{ } F \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-map-left*:

$\langle \text{kf-comp (kf-map } f \text{ } E) \text{ } F \equiv_{kr} \text{ kf-map } (\lambda(x,y). (x, f \text{ } y)) \text{ (kf-comp } E \text{ } F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-map-right*:

$\langle \text{kf-comp } E \text{ (kf-map } f \text{ } F) \equiv_{kr} \text{ kf-map } (\lambda(x,y). (f \text{ } x, y)) \text{ (kf-comp } E \text{ } F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-map-both*:

$\langle \text{kf-comp (kf-map } e \text{ } E) \text{ (kf-map } f \text{ } F) \equiv_{kr} \text{ kf-map } (\lambda(x,y). (f \text{ } x, e \text{ } y)) \text{ (kf-comp } E \text{ } F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-commute*:

assumes $\langle \text{kf-operators } \mathfrak{F} \subseteq \text{commutant (kf-operators } \mathfrak{E}) \rangle$
shows $\langle \text{kf-apply } \mathfrak{F} \circ \text{kf-apply } \mathfrak{E} = \text{kf-apply } \mathfrak{E} \circ \text{kf-apply } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-commute-weak*:

assumes $\langle \text{kf-operators } \mathfrak{F} \subseteq \text{commutant (kf-operators } \mathfrak{E}) \rangle$
shows $\langle \text{kf-comp } \mathfrak{F} \text{ } \mathfrak{E} \equiv_{kr} \text{kf-comp } \mathfrak{E} \text{ } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-commute*:

assumes $\langle \text{kf-operators } \mathfrak{F} \subseteq \text{commutant (kf-operators } \mathfrak{E}) \rangle$
shows $\langle \text{kf-comp } \mathfrak{F} \text{ } \mathfrak{E} \equiv_{kr} \text{kf-map prod.swap (kf-comp } \mathfrak{E} \text{ } \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-apply-on-singleton*:

$\langle \text{kf-comp } \mathfrak{E} \text{ } \mathfrak{F} *_{kr} @\{x\} \text{ } \varrho = \mathfrak{E} *_{kr} @\{\text{snd } x\} \text{ } (\mathfrak{F} *_{kr} @\{\text{fst } x\} \text{ } \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-apply-on-singleton*:

assumes $\langle \text{bdd-above ((kf-norm } \circ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \text{ } \mathfrak{F} *_{kr} @\{x\} \text{ } \varrho = \mathfrak{E} (\text{fst } x) *_{kr} @\{\text{snd } x\} \text{ } (\mathfrak{F} *_{kr} @\{\text{fst } x\} \text{ } \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-id-left*: $\langle \text{kf-comp kf-id } \mathfrak{E} \equiv_{kr} \text{kf-map } (\lambda x. (x, ())) \text{ } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-comp-id-right*: $\langle \text{kf-comp } \mathfrak{E} \text{ kf-id } \equiv_{kr} \text{kf-map } (\lambda x. ((), x)) \text{ } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-kf-plus-left*:

fixes $\mathfrak{D} :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'd) \text{ kraus-family} \rangle$
fixes $\mathfrak{E} :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$

fixes $\mathfrak{F} :: \langle ('a :: \text{hilbert-space}, 'b, 'f) \text{ kraus-family} \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{D} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{E} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent-raw } (\lambda x. \text{kf-plus } (\mathfrak{D} x) (\mathfrak{E} x)) \mathfrak{F} =$
 $\text{kf-map-inj } (\lambda x. \text{case } x \text{ of Inl } (F, D, f, d) \Rightarrow (F, D, f, \text{Inl } d) \mid \text{Inr } (F, E, f, e) \Rightarrow (F, E, f, \text{Inr } e))$
 $(\text{kf-plus } (\text{kf-comp-dependent-raw } \mathfrak{D} \mathfrak{F}) (\text{kf-comp-dependent-raw } \mathfrak{E} \mathfrak{F})) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-kf-plus-left*:
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{D} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{E} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } (\lambda x. \text{kf-plus } (\mathfrak{D} x) (\mathfrak{E} x)) \mathfrak{F} =$
 $\text{kf-map-inj } (\lambda x. \text{case } x \text{ of Inl } (f, d) \Rightarrow (f, \text{Inl } d) \mid \text{Inr } (f, e) \Rightarrow (f, \text{Inr } e)) (\text{kf-plus } (\text{kf-comp-dependent } \mathfrak{D} \mathfrak{F}) (\text{kf-comp-dependent } \mathfrak{E} \mathfrak{F})) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-twice*:
shows $\langle \text{kf-map-inj } f (\text{kf-map-inj } g \mathfrak{E}) = \text{kf-map-inj } (f \circ g) \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-kf-plus-left'*:
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{D} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{E} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-plus } (\text{kf-comp-dependent } \mathfrak{D} \mathfrak{F}) (\text{kf-comp-dependent } \mathfrak{E} \mathfrak{F}) =$
 $\text{kf-map-inj } (\lambda (f, de). \text{case } de \text{ of Inl } d \Rightarrow \text{Inl } (f, d) \mid \text{Inr } e \Rightarrow \text{Inr } (f, e)) (\text{kf-comp-dependent } (\lambda x. \text{kf-plus } (\mathfrak{D} x) (\mathfrak{E} x)) \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-plus-left*:
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{D} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{E} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } (\lambda x. \mathfrak{D} x + \mathfrak{E} x) \mathfrak{F} \equiv_{kr} \text{kf-comp-dependent } \mathfrak{D} \mathfrak{F} + \text{kf-comp-dependent } \mathfrak{E} \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

2.9 Infinite sums

lift-definition *kf-infsum* :: $\langle ('a \Rightarrow ('b :: \text{hilbert-space}, 'c :: \text{hilbert-space}, 'x) \text{ kraus-family}) \Rightarrow 'a \text{ set} \Rightarrow ('b, 'c, 'a \times 'x) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda \mathfrak{E} A. \text{if summable-on-in cweak-operator-topology } (\lambda a. \text{kf-bound } (\mathfrak{E} a)) A$
 $\text{then } (\lambda (a, (E, x)). (E, (a, x))) \text{ 'Sigma } A (\lambda a. \text{Rep-kraus-family } (\mathfrak{E} a)) \text{ else } \{\} \rangle$
 $\langle \text{proof} \rangle$

definition *kf-summable* :: $\langle ('a \Rightarrow ('b :: \text{hilbert-space}, 'c :: \text{hilbert-space}, 'x) \text{ kraus-family}) \Rightarrow 'a \text{ set} \Rightarrow \text{bool} \rangle$ **where**

$\langle \text{kf-summable } \mathfrak{E} A \longleftrightarrow \text{summable-on-in cweak-operator-topology } (\lambda a. \text{kf-bound } (\mathfrak{E} a)) A \rangle$

lemma *kf-summable-from-abs-summable*:

assumes *sum*: $\langle (\lambda a. \text{kf-norm } (\mathfrak{E} a)) \text{ summable-on } A \rangle$

shows $\langle \text{kf-summable } \mathfrak{E} A \rangle$

<proof>

lemma *kf-infsum-apply*:

assumes $\langle \text{kf-summable } \mathfrak{E} A \rangle$

shows $\langle \text{kf-infsum } \mathfrak{E} A *_{kr} \varrho = (\sum_{\infty} a \in A. \mathfrak{E} a *_{kr} \varrho) \rangle$

<proof>

lemma *kf-infsum-apply-summable*:

assumes $\langle \text{kf-summable } \mathfrak{E} A \rangle$

shows $\langle (\lambda a. \mathfrak{E} a *_{kr} \varrho) \text{ summable-on } A \rangle$

<proof>

lemma *kf-bound-infsum*:

fixes *f* :: $\langle 'a \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$

assumes $\langle \text{kf-summable } f A \rangle$

shows $\langle \text{kf-bound } (\text{kf-infsum } f A) = \text{infsum-in cweak-operator-topology } (\lambda a. \text{kf-bound } (f a)) A \rangle$

<proof>

lemma *kf-norm-infsum*:

assumes *sum*: $\langle (\lambda a. \text{kf-norm } (\mathfrak{E} a)) \text{ summable-on } A \rangle$

shows $\langle \text{kf-norm } (\text{kf-infsum } \mathfrak{E} A) \leq (\sum_{\infty} a \in A. \text{kf-norm } (\mathfrak{E} a)) \rangle$

<proof>

lemma *kf-filter-infsum*:

assumes $\langle \text{kf-summable } \mathfrak{E} A \rangle$

shows $\langle \text{kf-filter } P (\text{kf-infsum } \mathfrak{E} A)$

$= \text{kf-infsum } (\lambda a. \text{kf-filter } (\lambda x. P (a, x)) (\mathfrak{E} a)) \{a \in A. \exists x. P (a, x)\} \rangle$

(**is** $\langle ?lhs = ?rhs \rangle$)

<proof>

lemma *kf-infsum-empty[simp]*: $\langle \text{kf-infsum } \mathfrak{E} \{\} = 0 \rangle$

<proof>

lemma *kf-infsum-singleton[simp]*: $\langle \text{kf-infsum } \mathfrak{E} \{a\} = \text{kf-map-inj } (\lambda x. (a, x)) (\mathfrak{E} a) \rangle$

<proof>

lemma *kf-infsum-invalid*:

assumes $\langle \neg \text{kf-summable } \mathfrak{E} A \rangle$

shows $\langle \text{kf-infsum } \mathfrak{E} A = 0 \rangle$

<proof>

lemma *kf-infsum-cong*:

fixes $\mathfrak{E} \mathfrak{F} :: \langle 'a \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$

assumes $\langle \bigwedge a. a \in A \implies \mathfrak{E} a \equiv_{kr} \mathfrak{F} a \rangle$
shows $\langle \text{kf-infsum } \mathfrak{E} A \equiv_{kr} \text{kf-infsum } \mathfrak{F} A \rangle$
 $\langle \text{proof} \rangle$

2.10 Trace-preserving maps

definition $\langle \text{kf-trace-preserving } \mathfrak{E} \longleftrightarrow (\forall \varrho. \text{trace-tc } (\mathfrak{E} *_{kr} \varrho) = \text{trace-tc } \varrho) \rangle$

definition $\langle \text{kf-trace-reducing } \mathfrak{E} \longleftrightarrow (\forall \varrho \geq 0. \text{trace-tc } (\mathfrak{E} *_{kr} \varrho) \leq \text{trace-tc } \varrho) \rangle$

lemma *kf-trace-reducing-iff-norm-leq1*: $\langle \text{kf-trace-reducing } \mathfrak{E} \longleftrightarrow \text{kf-norm } \mathfrak{E} \leq 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-trace-preserving-iff-bound-id*: $\langle \text{kf-trace-preserving } \mathfrak{E} \longleftrightarrow \text{kf-bound } \mathfrak{E} = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-trace-norm-preserving*: $\langle \text{kf-norm } \mathfrak{E} \leq 1 \rangle$ **if** $\langle \text{kf-trace-preserving } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-trace-norm-preserving-eq*:
fixes $\mathfrak{E} :: \langle ('a :: \{\text{chilbert-space, not-singleton}\}, 'b :: \text{chilbert-space}, 'c) \text{ kraus-family} \rangle$
assumes $\langle \text{kf-trace-preserving } \mathfrak{E} \rangle$
shows $\langle \text{kf-norm } \mathfrak{E} = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-trace-preserving-map[simp]*: $\langle \text{kf-trace-preserving } (\text{kf-map } f \ \mathfrak{E}) \longleftrightarrow \text{kf-trace-preserving } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-trace-reducing-map[simp]*: $\langle \text{kf-trace-reducing } (\text{kf-map } f \ \mathfrak{E}) \longleftrightarrow \text{kf-trace-reducing } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-trace-preserving-id[iff]*: $\langle \text{kf-trace-preserving } \text{kf-id} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-trace-reducing-id[iff]*: $\langle \text{kf-trace-reducing } \text{kf-id} \rangle$
 $\langle \text{proof} \rangle$

2.11 Sampling

lift-definition *kf-sample* :: $\langle ('x \Rightarrow \text{real}) \Rightarrow ('a :: \text{chilbert-space}, 'a, 'x) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda p. \text{if } (\forall x. p \ x \geq 0) \wedge p \text{ summable-on } \text{UNIV} \text{ then } \text{Set.filter } (\lambda(E, -). E \neq 0) (\text{range } (\lambda x. (\text{sqrt } (p \ x) *_{kr} \text{id-cblinfun}, x))) \text{ else } \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-sample-norm*:
fixes $p :: \langle 'x \Rightarrow \text{real} \rangle$
assumes $\langle \bigwedge x. p \ x \geq 0 \rangle$
assumes $\langle p \text{ summable-on } \text{UNIV} \rangle$

shows $\langle \text{kf-norm } (\text{kf-sample } p :: ('a::\{\text{chilbert-space, not-singleton}\}, 'a, 'x) \text{ kraus-family})$
 $= (\sum_{\infty x. p \ x}) \rangle$
 $\langle \text{proof} \rangle$

2.12 Trace

lift-definition $\text{kf-trace} :: \langle 'a \text{ set} \Rightarrow ('a::\text{chilbert-space}, 'b::\text{one-dim}, 'a) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda B. \text{ if is-onb } B \text{ then } (\lambda x. (\text{vector-to-cblinfun } x*, x)) \text{ ' } B \text{ else } \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma kf-trace-is-trace :
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-trace } B *_{kr} \varrho = \text{one-dim-iso } (\text{trace-tc } \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{kf-eq-weak-kf-trace}$:
assumes $\langle \text{is-onb } A \rangle$ **and** $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-trace } A =_{kr} \text{kf-trace } B \rangle$
 $\langle \text{proof} \rangle$

lemma trace-is-kf-trace :
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{trace-tc } t = \text{one-dim-iso } (\text{kf-trace } B *_{kr} t) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{kf-trace-bound[simp]}$:
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-bound } (\text{kf-trace } B) = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{kf-trace-norm-eq1[simp]}$:
fixes $B :: \langle 'a::\{\text{chilbert-space, not-singleton}\} \text{ set} \rangle$
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-norm } (\text{kf-trace } B) = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{kf-trace-norm-leq1[simp]}$:
fixes $B :: \langle 'a::\text{chilbert-space set} \rangle$
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-norm } (\text{kf-trace } B) \leq 1 \rangle$
 $\langle \text{proof} \rangle$

2.13 Constant maps

lift-definition $\text{kf-constant-onedim} :: \langle ('b, 'b) \text{ trace-class} \Rightarrow ('a::\text{one-dim}, 'b::\text{chilbert-space}, \text{unit})$
 $\text{kraus-family} \rangle$ **is**
 $\langle \lambda t::('b, 'b) \text{ trace-class. if } t \geq 0 \text{ then}$
 $\text{Set.filter } (\lambda(E, -). E \neq 0) ((\lambda v. (\text{vector-to-cblinfun } v, ())) \text{ ' spectral-dec-vecs-tc } t)$
 $\text{else } \{\} \rangle$

⟨proof⟩

definition $\text{kf-constant} :: \langle 'b, 'b \rangle \text{ trace-class} \Rightarrow \langle 'a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, \text{unit} \rangle \text{ kraus-family}$
where

$\langle \text{kf-constant } \varrho = \text{kf-flatten } (\text{kf-comp } (\text{kf-constant-onedim } \varrho :: (\text{complex}, -, -) \text{ kraus-family}) (\text{kf-trace } \text{some-chilbert-basis})) \rangle$

lemma $\text{kf-constant-onedim-invalid}: \langle \neg \varrho \geq 0 \implies \text{kf-constant-onedim } \varrho = 0 \rangle$
 ⟨proof⟩

lemma $\text{kf-constant-invalid}: \langle \neg \varrho \geq 0 \implies \text{kf-constant } \varrho = 0 \rangle$
 ⟨proof⟩

lemma $\text{kf-constant-onedim-apply}$:
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-apply } (\text{kf-constant-onedim } \varrho) \sigma = \text{one-dim-iso } \sigma *_C \varrho \rangle$
 ⟨proof⟩

lemma kf-constant-apply :
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-apply } (\text{kf-constant } \varrho) \sigma = \text{trace-tc } \sigma *_C \varrho \rangle$
 ⟨proof⟩

lemma $\text{kf-bound-constant-onedim[simp]}$:
fixes $\varrho :: \langle 'a :: \text{chilbert-space}, 'a \rangle \text{ trace-class}$
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-bound } (\text{kf-constant-onedim } \varrho) = \text{norm } \varrho *_R \text{id-cblinfun} \rangle$
 ⟨proof⟩

lemma $\text{kf-bound-constant[simp]}$:
fixes $\varrho :: \langle 'a :: \text{chilbert-space}, 'a \rangle \text{ trace-class}$
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-bound } (\text{kf-constant } \varrho) = \text{norm } \varrho *_R \text{id-cblinfun} \rangle$
 ⟨proof⟩

lemma $\text{kf-norm-constant-onedim[simp]}$:
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-norm } (\text{kf-constant-onedim } \varrho) = \text{norm } \varrho \rangle$
 ⟨proof⟩

lemma kf-norm-constant :
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-norm } (\text{kf-constant } \varrho :: \langle 'a :: \{ \text{chilbert-space}, \text{not-singleton} \}, 'b :: \text{chilbert-space}, - \rangle \text{ kraus-family}) = \text{norm } \varrho \rangle$
 ⟨proof⟩

lemma $\text{kf-norm-constant-leq}$:
shows $\langle \text{kf-norm } (\text{kf-constant } \varrho) \leq \text{norm } \varrho \rangle$
 ⟨proof⟩

lemma *kf-comp-constant-right*:

assumes $[iff]: \langle t \geq 0 \rangle$

shows $\langle kf\text{-map } fst \ (kf\text{-comp } E \ (kf\text{-constant } t)) \equiv_{kr} kf\text{-constant } (E *_{kr} t) \rangle$

$\langle proof \rangle$

lemma *kf-comp-constant-right-weak*:

assumes $[iff]: \langle t \geq 0 \rangle$

shows $\langle kf\text{-comp } E \ (kf\text{-constant } t) =_{kr} kf\text{-constant } (E *_{kr} t) \rangle$

$\langle proof \rangle$

2.14 Tensor products

lemma *kf-tensor-raw-bound-aux*:

fixes $\mathfrak{E} :: \langle ('a \text{ ell2} \Rightarrow_{CL} 'b \text{ ell2} \times 'x) \text{ set} \rangle$ **and** $\mathfrak{F} :: \langle ('c \text{ ell2} \Rightarrow_{CL} 'd \text{ ell2} \times 'y) \text{ set} \rangle$

assumes $\langle \bigwedge S. \text{finite } S \implies S \subseteq \mathfrak{E} \implies (\sum_{(E,x) \in S} E * o_{CL} E) \leq B \rangle$

assumes $\langle \bigwedge S. \text{finite } S \implies S \subseteq \mathfrak{F} \implies (\sum_{(E,x) \in S} E * o_{CL} E) \leq C \rangle$

assumes $\langle \text{finite } U \rangle$

assumes $\langle U \subseteq ((\lambda((E,x), (F,y)). (E \otimes_o F, E, F, x, y))) \text{ ' } (\mathfrak{E} \times \mathfrak{F}) \rangle$

shows $\langle (\sum_{(G,z) \in U} G * o_{CL} G) \leq B \otimes_o C \rangle$

$\langle proof \rangle$

lift-definition *kf-tensor-raw* :: $\langle ('a \text{ ell2}, 'b \text{ ell2}, 'x) \text{ kraus-family} \Rightarrow ('c \text{ ell2}, 'd \text{ ell2}, 'y) \text{ kraus-family} \Rightarrow$

$((('a \times 'c) \text{ ell2}, ('b \times 'd) \text{ ell2}, ((('a \text{ ell2} \Rightarrow_{CL} 'b \text{ ell2}) \times ('c \text{ ell2} \Rightarrow_{CL} 'd \text{ ell2}) \times 'x \times 'y)) \text{ kraus-family} \rangle$

is

$\langle \lambda \mathfrak{E} \mathfrak{F}. (\lambda((E,x), (F,y)). (E \otimes_o F, (E, F, x, y))) \text{ ' } (\mathfrak{E} \times \mathfrak{F}) \rangle$

$\langle proof \rangle$

lemma *kf-apply-tensor-raw-as-infsum*:

$\langle kf\text{-tensor-raw } \mathfrak{E} \mathfrak{F} *_{kr} \varrho = (\sum_{\infty} ((E,x), (F,y)) \in \text{Rep-kraus-family } \mathfrak{E} \times \text{Rep-kraus-family } \mathfrak{F}. \text{ sandwich-tc } (E \otimes_o F) \varrho) \rangle$

$\langle proof \rangle$

lemma *kf-apply-tensor-raw*:

shows $\langle kf\text{-tensor-raw } \mathfrak{E} \mathfrak{F} *_{kr} \text{tc-tensor } \varrho \sigma = \text{tc-tensor } (\mathfrak{E} *_{kr} \varrho) (\mathfrak{F} *_{kr} \sigma) \rangle$

$\langle proof \rangle$

hide-fact *kf-tensor-raw-bound-aux*

definition $\langle kf\text{-tensor } \mathfrak{E} \mathfrak{F} = kf\text{-map } (\lambda(E, F, x, y). (x, y)) \ (kf\text{-tensor-raw } \mathfrak{E} \mathfrak{F}) \rangle$

lemma *kf-apply-tensor*:

$\langle kf\text{-tensor } \mathfrak{E} \mathfrak{F} *_{kr} \text{tc-tensor } \varrho \sigma = \text{tc-tensor } (\mathfrak{E} *_{kr} \varrho) (\mathfrak{F} *_{kr} \sigma) \rangle$

$\langle proof \rangle$

lemma *kf-apply-tensor-as-infsum*:

$\langle kf\text{-tensor } \mathfrak{E} \mathfrak{F} *_{kr} \varrho = (\sum_{\infty} ((E,x), (F,y)) \in \text{Rep-kraus-family } \mathfrak{E} \times \text{Rep-kraus-family } \mathfrak{F}. \text{ sand-}$

wich-tc ($E \otimes_o F$) ϱ ›
 ⟨*proof*⟩

lemma *kf-bound-tensor-raw*:

⟨*kf-bound* (*kf-tensor-raw* $\mathfrak{E}$ $\mathfrak{F}$) = *kf-bound* $\mathfrak{E} \otimes_o$ *kf-bound* $\mathfrak{F}$ ›
 ⟨*proof*⟩

lemma *kf-bound-tensor*:

⟨*kf-bound* (*kf-tensor* $\mathfrak{E}$ $\mathfrak{F}$) = *kf-bound* $\mathfrak{E} \otimes_o$ *kf-bound* $\mathfrak{F}$ ›
 ⟨*proof*⟩

lemma *kf-norm-tensor*:

⟨*kf-norm* (*kf-tensor* $\mathfrak{E}$ $\mathfrak{F}$) = *kf-norm* $\mathfrak{E} *$ *kf-norm* $\mathfrak{F}$ ›
 ⟨*proof*⟩

lemma *kf-tensor-cong-weak*:

assumes ⟨ $\mathfrak{E} =_{kr} \mathfrak{E}'$ ›
assumes ⟨ $\mathfrak{F} =_{kr} \mathfrak{F}'$ ›
shows ⟨*kf-tensor* $\mathfrak{E}$ $\mathfrak{F} =_{kr}$ *kf-tensor* $\mathfrak{E}'$ $\mathfrak{F}'$ ›
 ⟨*proof*⟩

lemma *kf-filter-tensor*:

⟨*kf-filter* ($\lambda(x,y). P\ x \wedge Q\ y$) (*kf-tensor* $\mathfrak{E}$ $\mathfrak{F}$) = *kf-tensor* (*kf-filter* P $\mathfrak{E}$) (*kf-filter* Q $\mathfrak{F}$)›
 ⟨*proof*⟩

lemma *kf-filter-tensor-singleton*:

⟨*kf-filter* ($((=)\ x)$) (*kf-tensor* $\mathfrak{E}$ $\mathfrak{F}$) = *kf-tensor* (*kf-filter* ($((=)\ (fst\ x))$) $\mathfrak{E}$) (*kf-filter* ($((=)\ (snd\ x))$) $\mathfrak{F}$)›
 ⟨*proof*⟩

lemma *kf-tensor-cog*:

fixes $\mathfrak{E}\ \mathfrak{E}' :: \langle ('a\ ell2, 'b\ ell2, 'x)\ kraus-family \rangle$
and $\mathfrak{F}\ \mathfrak{F}' :: \langle ('c\ ell2, 'd\ ell2, 'y)\ kraus-family \rangle$
assumes ⟨ $\mathfrak{E} \equiv_{kr} \mathfrak{E}'$ ›
assumes ⟨ $\mathfrak{F} \equiv_{kr} \mathfrak{F}'$ ›
shows ⟨*kf-tensor* $\mathfrak{E}$ $\mathfrak{F} \equiv_{kr}$ *kf-tensor* $\mathfrak{E}'$ $\mathfrak{F}'$ ›
 ⟨*proof*⟩

lemma *kf-tensor-compose-distrib-weak*:

shows ⟨*kf-tensor* (*kf-comp* $\mathfrak{E}$ $\mathfrak{F}$) (*kf-comp* $\mathfrak{G}$ $\mathfrak{H}$)
 =_{kr} *kf-comp* (*kf-tensor* $\mathfrak{E}$ $\mathfrak{G}$) (*kf-tensor* $\mathfrak{F}$ $\mathfrak{H}$)›
 ⟨*proof*⟩

lemma *kf-tensor-compose-distrib*:

shows ⟨*kf-tensor* (*kf-comp* $\mathfrak{E}$ $\mathfrak{F}$) (*kf-comp* $\mathfrak{G}$ $\mathfrak{H}$)
 $\equiv_{kr}$ *kf-map* ($\lambda((e,g),(f,h)). ((e,f),(g,h))$) (*kf-comp* (*kf-tensor* $\mathfrak{E}$ $\mathfrak{G}$) (*kf-tensor* $\mathfrak{F}$ $\mathfrak{H}$))›

⟨proof⟩

lemma *kf-tensor-compose-distrib'*:

shows ⟨*kf-comp* (*kf-tensor* $\mathfrak{E}$ $\mathfrak{G}$) (*kf-tensor* $\mathfrak{F}$ $\mathfrak{H}$)

$\equiv_{kr}$ *kf-map* ($\lambda((e,f),(g,h)). ((e,g),(f,h)))$ (*kf-tensor* (*kf-comp* $\mathfrak{E}$ $\mathfrak{F}$) (*kf-comp* $\mathfrak{G}$ $\mathfrak{H}$)))⟩

⟨proof⟩

definition *kf-tensor-right* :: ⟨('extra ell2, 'extra ell2) trace-class $\Rightarrow$ ('qu ell2, ('qu $\times$ 'extra) ell2, unit) kraus-family⟩ **where**

— *kf-tensor-right* ϱ maps σ to $\sigma \otimes_o \varrho$

⟨*kf-tensor-right* ϱ = *kf-map-inj* ($\lambda. ()$) (*kf-comp* (*kf-tensor* *kf-id* (*kf-constant-onedim* ϱ)) (*kf-of-op* (*tensor-ell2-right* (*ket* ())))))⟩

definition *kf-tensor-left* :: ⟨('extra ell2, 'extra ell2) trace-class $\Rightarrow$ ('qu ell2, ('extra $\times$ 'qu) ell2, unit) kraus-family⟩ **where**

— *kf-tensor-right* ϱ maps σ to $\varrho \otimes_o \sigma$

⟨*kf-tensor-left* ϱ = *kf-map-inj* ($\lambda. ()$) (*kf-comp* (*kf-tensor* (*kf-constant-onedim* ϱ) *kf-id*) (*kf-of-op* (*tensor-ell2-left* (*ket* ())))))⟩

lemma *kf-apply-tensor-right[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-tensor-right* ϱ *_{kr} σ = *tc-tensor* σ ϱ ⟩

⟨proof⟩

lemma *kf-apply-tensor-left[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-tensor-left* ϱ *_{kr} σ = *tc-tensor* ϱ σ ⟩

⟨proof⟩

lemma *kf-bound-tensor-right[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-bound* (*kf-tensor-right* ϱ) = *norm* ϱ *_C *id-cblinfun*⟩

⟨proof⟩

lemma *kf-bound-tensor-left[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-bound* (*kf-tensor-left* ϱ) = *norm* ϱ *_C *id-cblinfun*⟩

⟨proof⟩

lemma *kf-norm-tensor-right[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-norm* (*kf-tensor-right* ϱ) = *norm* ϱ ⟩

⟨proof⟩

lemma *kf-norm-tensor-left[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-norm* (*kf-tensor-left* ϱ) = *norm* ϱ ⟩

⟨proof⟩

lemma *kf-trace-preserving-tensor*:

assumes $\langle \text{kf-trace-preserving } \mathfrak{E} \rangle$ **and** $\langle \text{kf-trace-preserving } \mathfrak{F} \rangle$
shows $\langle \text{kf-trace-preserving } (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-trace-reducing-tensor*:
assumes $\langle \text{kf-trace-reducing } \mathfrak{E} \rangle$ **and** $\langle \text{kf-trace-reducing } \mathfrak{F} \rangle$
shows $\langle \text{kf-trace-reducing } (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-map-left*:
 $\langle \text{kf-tensor } (\text{kf-map } f \ \mathfrak{E}) \ \mathfrak{F} \equiv_{kr} \text{kf-map } (\text{apfst } f) (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-map-right*:
 $\langle \text{kf-tensor } \mathfrak{E} (\text{kf-map } f \ \mathfrak{F}) \equiv_{kr} \text{kf-map } (\text{apsnd } f) (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-map-both*:
 $\langle \text{kf-tensor } (\text{kf-map } f \ \mathfrak{E}) (\text{kf-map } g \ \mathfrak{F}) \equiv_{kr} \text{kf-map } (\text{map-prod } f \ g) (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-raw-map-inj-both*:
 $\langle \text{kf-tensor-raw } (\text{kf-map-inj } f \ \mathfrak{E}) (\text{kf-map-inj } g \ \mathfrak{F}) = \text{kf-map-inj } (\lambda(E, F, x, y). (E, F, f \ x, g \ y)) (\text{kf-tensor-raw } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-tensor-raw-subset*:
 $\langle \text{kf-domain } (\text{kf-tensor-raw } \mathfrak{E} \ \mathfrak{F}) \subseteq \text{kf-operators } \mathfrak{E} \times \text{kf-operators } \mathfrak{F} \times \text{kf-domain } \mathfrak{E} \times \text{kf-domain } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-map-inj-both*:
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
assumes $\langle \text{inj-on } g \ (\text{kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-tensor } (\text{kf-map-inj } f \ \mathfrak{E}) (\text{kf-map-inj } g \ \mathfrak{F}) = \text{kf-map-inj } (\text{map-prod } f \ g) (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-tensor-raw*:
shows $\langle \text{kf-operators } (\text{kf-tensor-raw } \mathfrak{E} \ \mathfrak{F}) = \{E \otimes_o F \mid E \ F. E \in \text{kf-operators } \mathfrak{E} \wedge F \in \text{kf-operators } \mathfrak{F}\} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-tensor*:
shows $\langle \text{kf-operators } (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \subseteq \text{span } \{E \otimes_o F \mid E \ F. E \in \text{kf-operators } \mathfrak{E} \wedge F \in \text{kf-operators } \mathfrak{F}\} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-tensor*: $\langle \text{kf-domain } (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) = \text{kf-domain } \mathfrak{E} \times \text{kf-domain } \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-tensor-raw-0-left[simp]*: $\langle \text{kf-tensor-raw } 0 \ \mathfrak{E} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-raw-0-right[simp]*: $\langle \text{kf-tensor-raw } \mathfrak{E} \ 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-0-left[simp]*: $\langle \text{kf-tensor } 0 \ \mathfrak{E} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-0-right[simp]*: $\langle \text{kf-tensor } \mathfrak{E} \ 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-of-op*:
 $\langle \text{kf-tensor } (\text{kf-of-op } A) (\text{kf-of-op } B) = \text{kf-map } (\lambda(). ((), ())) (\text{kf-of-op } (A \otimes_o B)) \rangle$
 $\langle \text{proof} \rangle$

2.15 Partial trace

definition *kf-partial-trace-right* :: $\langle (('a \times 'b) \ \text{ell2}, 'a \ \text{ell2}, 'b) \ \text{kraus-family} \rangle$ **where**
 $\langle \text{kf-partial-trace-right} = \text{kf-map } (\lambda(-, b). \text{inv ket } b)$
 $(\text{kf-comp } (\text{kf-of-op } (\text{tensor-ell2-right } (\text{ket } ())) *))$
 $(\text{kf-tensor kf-id } (\text{kf-trace } (\text{range ket})))) \rangle$

definition *kf-partial-trace-left* :: $\langle (('a \times 'b) \ \text{ell2}, 'b \ \text{ell2}, 'a) \ \text{kraus-family} \rangle$ **where**
 $\langle \text{kf-partial-trace-left} = \text{kf-map-inj snd } (\text{kf-comp } \text{kf-partial-trace-right } (\text{kf-of-op swap-ell2})) \rangle$

lemma *partial-trace-is-kf-partial-trace*:
fixes $t :: \langle (('a \times 'b) \ \text{ell2}, ('a \times 'b) \ \text{ell2}) \ \text{trace-class} \rangle$
shows $\langle \text{partial-trace } t = \text{kf-partial-trace-right } *_{kr} \ t \rangle$
 $\langle \text{proof} \rangle$

lemma *partial-trace-ignores-kraus-family*:
assumes $\langle \text{kf-trace-preserving } \mathfrak{E} \rangle$
shows $\langle \text{partial-trace } (\text{kf-tensor } \mathfrak{F} \ \mathfrak{E} \ *_{kr} \ \varrho) = \mathfrak{F} \ *_{kr} \ \text{partial-trace } \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-partial-trace-bound[simp]*:
shows $\langle \text{kf-bound } \text{kf-partial-trace-right} = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-partial-trace-norm[simp]*:
shows $\langle \text{kf-norm } \text{kf-partial-trace-right} = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-partial-trace-right-apply-singleton*:
 $\langle \text{kf-partial-trace-right } *_{kr} \ @\{x\} \ \varrho = \text{sandwich-tc } (\text{tensor-ell2-right } (\text{ket } x) *) \ \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-partial-trace-left-apply-singleton*:

$\langle \text{kf-partial-trace-left } *_{kr} \text{ @}\{x\} \varrho = \text{sandwich-tc } (\text{tensor-ell2-left } (\text{ket } x) *) \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-partial-trace-right[simp]*: $\langle \text{kf-domain } \text{kf-partial-trace-right} = \text{UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-partial-trace-left[simp]*: $\langle \text{kf-domain } \text{kf-partial-trace-left} = \text{UNIV} \rangle$
 $\langle \text{proof} \rangle$

2.16 Complete measurement

lemma *complete-measurement-aux*:

fixes B **and** $F :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'a \times 'a \text{ set} \rangle$
defines $\langle \text{family} \equiv (\lambda x. (\text{selfbutter } (\text{sgn } x), x)) \text{ ' } B \rangle$
assumes $\langle \text{finite } F \rangle$ **and** $FB: \langle F \subseteq \text{family} \rangle$ **and** $\langle \text{is-ortho-set } B \rangle$
shows $\langle (\sum (E, x) \in F. E * o_{CL} E) \leq \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *complete-measurement-is-kraus-family*:

assumes $\langle \text{is-ortho-set } B \rangle$
shows $\langle \text{kraus-family } ((\lambda x. (\text{selfbutter } (\text{sgn } x), x)) \text{ ' } B) \rangle$
 $\langle \text{proof} \rangle$

lift-definition *kf-complete-measurement* :: $\langle 'a \text{ set} \Rightarrow ('a :: \text{chilbert-space}, 'a, 'a) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda B. \text{if is-ortho-set } B \text{ then } (\lambda x. (\text{selfbutter } (\text{sgn } x), x)) \text{ ' } B \text{ else } \{\} \rangle$
 $\langle \text{proof} \rangle$

definition *kf-complete-measurement-ket* :: $\langle ('a \text{ ell2}, 'a \text{ ell2}, 'a) \text{ kraus-family} \rangle$ **where**
 $\langle \text{kf-complete-measurement-ket} = \text{kf-map-inj } (\text{inv ket}) (\text{kf-complete-measurement } (\text{range ket})) \rangle$

lemma *kf-complete-measurement-domain[simp]*:

assumes $\langle \text{is-ortho-set } B \rangle$
shows $\langle \text{kf-domain } (\text{kf-complete-measurement } B) = B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-domain[simp]*:

$\langle \text{kf-domain } \text{kf-complete-measurement-ket} = \text{UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-kf-map*:

$\langle \text{kf-complete-measurement-ket} \equiv_{kr} \text{kf-map } (\text{inv ket}) (\text{kf-complete-measurement } (\text{range ket})) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-complete-measurement*:

assumes $\langle is\text{-ortho-set } B \rangle$
shows $\langle kf\text{-bound } (kf\text{-complete-measurement } B) \leq id\text{-cblinfun} \rangle$
 $\langle proof \rangle$

lemma *kf-norm-complete-measurement*:
assumes $\langle is\text{-ortho-set } B \rangle$
shows $\langle kf\text{-norm } (kf\text{-complete-measurement } B) \leq 1 \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-invalid*:
assumes $\langle \neg is\text{-ortho-set } B \rangle$
shows $\langle kf\text{-complete-measurement } B = 0 \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-idem*:
 $\langle kf\text{-comp } (kf\text{-complete-measurement } B) (kf\text{-complete-measurement } B)$
 $\equiv_{kr} kf\text{-map } (\lambda b. (b, b)) (kf\text{-complete-measurement } B) \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-idem-weak*:
 $\langle kf\text{-comp } (kf\text{-complete-measurement } B) (kf\text{-complete-measurement } B)$
 $=_{kr} kf\text{-complete-measurement } B \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-ket-idem*:
 $\langle kf\text{-comp } kf\text{-complete-measurement-ket } kf\text{-complete-measurement-ket}$
 $\equiv_{kr} kf\text{-map } (\lambda b. (b, b)) kf\text{-complete-measurement-ket} \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-ket-idem-weak*:
 $\langle kf\text{-comp } kf\text{-complete-measurement-ket } kf\text{-complete-measurement-ket}$
 $=_{kr} kf\text{-complete-measurement-ket} \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-apply*:
assumes $[simp]: \langle is\text{-ortho-set } B \rangle$
shows $\langle kf\text{-complete-measurement } B *_{kr} t = (\sum_{x \in B} sandwich\text{-tc } (selfbutter (sgn x)) t) \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-has-sum*:
assumes $\langle is\text{-ortho-set } B \rangle$
shows $\langle ((\lambda x. sandwich\text{-tc } (selfbutter (sgn x)) \varrho) \text{ has-sum } kf\text{-complete-measurement } B *_{kr} \varrho) B \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-has-sum-onb*:
assumes $\langle is\text{-onb } B \rangle$

shows $\langle ((\lambda x. \text{ sandwich-tc } (\text{selfbutter } x) \varrho) \text{ has-sum kf-complete-measurement } B *_{kr} \varrho) B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-has-sum*:
 $\langle ((\lambda x. \text{ sandwich-tc } (\text{selfbutter } (\text{ket } x)) \varrho) \text{ has-sum kf-complete-measurement-ket } *_{kr} \varrho) \text{ UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-apply-onb*:
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-complete-measurement } B *_{kr} t = (\sum_{\infty x \in B. \text{ sandwich-tc } (\text{selfbutter } x) t) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-apply*: $\langle \text{kf-complete-measurement-ket } *_{kr} t = (\sum_{\infty x. \text{ sandwich-tc } (\text{selfbutter } (\text{ket } x)) t) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-complete-measurement-onb[simp]*:
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-bound } (\text{kf-complete-measurement } B) = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-complete-measurement-ket[simp]*:
 $\langle \text{kf-bound kf-complete-measurement-ket} = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-complete-measurement-onb[simp]*:
fixes $B :: \langle 'a :: \{\text{not-singleton, hilbert-space}\} \text{ set} \rangle$
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-norm } (\text{kf-complete-measurement } B) = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-complete-measurement-ket[simp]*:
 $\langle \text{kf-norm kf-complete-measurement-ket} = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-diagonal-operator[simp]*:
 $\langle \text{kf-complete-measurement-ket } *_{kr} \text{ diagonal-operator-tc } f = \text{diagonal-operator-tc } f \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-complete-measurement*:
 $\langle \text{kf-operators } (\text{kf-complete-measurement } B) = (\text{selfbutter } o \text{ sgn}) \text{ ` } B \rangle$ **if** $\langle \text{is-ortho-set } B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-complete-measurement-invalid*:
 $\langle \text{kf-operators } (\text{kf-complete-measurement } B) = \{\} \rangle$ **if** $\langle \neg \text{is-ortho-set } B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-complete-measurement-ket*:

⟨*kf-operators* *kf-complete-measurement-ket* = *range* ($\lambda c.$ *butterfly* (*ket* *c*) (*ket* *c*))⟩
 ⟨*proof*⟩

lemma *kf-complete-measurement-apply-butterfly*:

assumes ⟨*is-ortho-set* *B*⟩ **and** ⟨*b* ∈ *B*⟩

shows ⟨*kf-complete-measurement* *B* *_{kr} *tc-butterfly* *b* *b* = *tc-butterfly* *b* *b*⟩

⟨*proof*⟩

lemma *kf-complete-measurement-ket-apply-butterfly*:

⟨*kf-complete-measurement-ket* *_{kr} *tc-butterfly* (*ket* *x*) (*ket* *x*) = *tc-butterfly* (*ket* *x*) (*ket* *x*)⟩

⟨*proof*⟩

lemma *kf-map-eq-kf-map-inj-singleton*:

assumes ⟨*card-le-1* (*Rep-kraus-family* $\mathfrak{E}$)⟩

shows ⟨*kf-map* *f* $\mathfrak{E}$ = *kf-map-inj* *f* $\mathfrak{E}$ ⟩

⟨*proof*⟩

lemma *kf-map-eq-kf-map-inj-singleton'*:

assumes ⟨ $\bigwedge y.$ *card-le-1* (*Rep-kraus-family* (*kf-filter* ((=) *y*) $\mathfrak{E}$))⟩

assumes ⟨*inj-on* *f* (*kf-domain* $\mathfrak{E}$)⟩

shows ⟨*kf-map* *f* $\mathfrak{E}$ = *kf-map-inj* *f* $\mathfrak{E}$ ⟩

⟨*proof*⟩

lemma *kf-filter-singleton-kf-complete-measurement*:

assumes ⟨*x* ∈ *B*⟩ **and** ⟨*is-ortho-set* *B*⟩

shows ⟨*kf-filter* ((=) *x*) (*kf-complete-measurement* *B*) = *kf-map-inj* ($\lambda.$ *x*) (*kf-of-op* (*selfbutter* (*sgn* *x*)))⟩

⟨*proof*⟩

lemma *kf-filter-singleton-kf-complete-measurement'*:

assumes ⟨*x* ∈ *B*⟩ **and** ⟨*is-ortho-set* *B*⟩

shows ⟨*kf-filter* ((=) *x*) (*kf-complete-measurement* *B*) = *kf-map* ($\lambda.$ *x*) (*kf-of-op* (*selfbutter* (*sgn* *x*)))⟩

⟨*proof*⟩

lemma *kf-filter-disjoint*:

assumes ⟨ $\bigwedge x.$ *x* ∈ *kf-domain* $\mathfrak{E} \implies P$ *x* = False⟩

shows ⟨*kf-filter* *P* $\mathfrak{E}$ = 0⟩

⟨*proof*⟩

lemma *kf-complete-measurement-tensor*:

assumes $\langle \text{is-ortho-set } B \rangle$ **and** $\langle \text{is-ortho-set } C \rangle$
shows $\langle \text{kf-map } (\lambda(b,c). b \otimes_s c) (\text{kf-tensor } (\text{kf-complete-measurement } B) (\text{kf-complete-measurement } C)) \rangle$
 $= \text{kf-complete-measurement } ((\lambda(b,c). b \otimes_s c) \text{ ‘ } (B \times C)) \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-filter*: $\langle \text{card-le-1 } (\text{Rep-kraus-family } (\text{kf-filter } P \mathfrak{E})) \rangle$ **if** $\langle \text{card-le-1 } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-map-inj*[*iff*]: $\langle \text{card-le-1 } (\text{Rep-kraus-family } (\text{kf-map-inj } f \mathfrak{E})) \rangle$ **if** $\langle \text{card-le-1 } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-map*[*iff*]: $\langle \text{card-le-1 } (\text{Rep-kraus-family } (\text{kf-map } f \mathfrak{E})) \rangle$ **if** $\langle \text{card-le-1 } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-tensor-raw*[*iff*]: $\langle \text{card-le-1 } (\text{Rep-kraus-family } (\text{kf-tensor-raw } \mathfrak{E} \mathfrak{F})) \rangle$ **if** $\langle \text{card-le-1 } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$ **and** $\langle \text{card-le-1 } (\text{Rep-kraus-family } \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-tensor*[*iff*]: $\langle \text{card-le-1 } (\text{Rep-kraus-family } (\text{kf-tensor } \mathfrak{E} \mathfrak{F})) \rangle$ **if** $\langle \text{card-le-1 } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$ **and** $\langle \text{card-le-1 } (\text{Rep-kraus-family } \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-filter-complete-measurement*: $\langle \text{card-le-1 } (\text{Rep-kraus-family } (\text{kf-filter } ((=)x) (\text{kf-complete-measurement } B))) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-tensor*:
shows $\langle \text{kf-tensor } (\text{kf-complete-measurement-ket} :: (-, -, 'a) \text{ kraus-family}) (\text{kf-complete-measurement-ket} :: (-, -, 'b) \text{ kraus-family}) \rangle$
 $= \text{kf-complete-measurement-ket} \rangle$
 $\langle \text{proof} \rangle$

2.17 Reconstruction

lemma *kf-reconstruction-is-bounded-clinear*:
assumes $\langle \bigwedge \varrho. ((\lambda a. \text{sandwich-tc } (f \ a) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ A \rangle$
shows $\langle \text{bounded-clinear } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-reconstruction-is-kraus-family*:
assumes $\text{sum: } \langle \bigwedge \varrho. ((\lambda a. \text{sandwich-tc } (f \ a) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ A \rangle$
defines $\langle F \equiv \text{Set.filter } (\lambda(E, -). E \neq 0) ((\lambda a. (f \ a, \ a)) \text{ ‘ } A) \rangle$
shows $\langle \text{kraus-family } F \rangle$

⟨proof⟩

lemma *kf-reconstruction*:

assumes *sum*: $\langle \bigwedge \varrho. ((\lambda a. \text{ sandwich-tc } (f \ a) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ A \rangle$

defines $\langle F \equiv \text{Abs-kraus-family } (\text{Set.filter } (\lambda(E,-). E \neq 0) ((\lambda a. (f \ a, \ a)) \ 'A)) \rangle$

shows $\langle \text{kf-apply } F = \mathfrak{E} \rangle$

⟨proof⟩

2.18 Cleanup

unbundle *no cblinfun-syntax*

unbundle *no kraus-map-syntax*

end

3 Kraus maps

theory *Kraus-Maps*

imports *Kraus-Families*

begin

3.1 Kraus maps

unbundle *kraus-map-syntax*

unbundle *cblinfun-syntax*

definition *kraus-map* :: $\langle ('a::\text{chilbert-space}, 'a) \text{ trace-class} \Rightarrow ('b::\text{chilbert-space}, 'b) \text{ trace-class} \rangle$

$\Rightarrow \text{bool}$ **where**

kraus-map-def-raw: $\langle \text{kraus-map } \mathfrak{E} \longleftrightarrow (\exists EE :: ('a, 'b, \text{unit}) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } EE) \rangle$

lemma *kraus-map-def*: $\langle \text{kraus-map } \mathfrak{E} \longleftrightarrow (\exists EE :: ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } EE) \rangle$

— Has a more general type than the original definition

⟨proof⟩

lemma *kraus-mapI*:

assumes $\langle \mathfrak{E} = \text{kf-apply } \mathfrak{E}' \rangle$

shows $\langle \text{kraus-map } \mathfrak{E} \rangle$

⟨proof⟩

lemma *kraus-map-bounded-clinear*:

$\langle \text{bounded-clinear } \mathfrak{E} \rangle$ **if** $\langle \text{kraus-map } \mathfrak{E} \rangle$

⟨proof⟩

lemma *kraus-map-pos*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \varrho \geq 0 \rangle$

shows $\langle \mathfrak{E} \ \varrho \geq 0 \rangle$

⟨proof⟩

lemma *kraus-map-mono*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \varrho \geq \tau \rangle$
shows $\langle \mathfrak{E} \varrho \geq \mathfrak{E} \tau \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-kf-apply*[*iff*]: $\langle \text{kraus-map } (\text{kf-apply } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

definition *km-some-kraus-family* :: $\langle ((\text{'a}::\text{chilbert-space}, \text{'a}) \text{ trace-class} \Rightarrow (\text{'b}::\text{chilbert-space}, \text{'b}) \text{ trace-class}) \Rightarrow (\text{'a}, \text{'b}, \text{unit}) \text{ kraus-family} \rangle$ **where**
 $\langle \text{km-some-kraus-family } \mathfrak{E} = (\text{if kraus-map } \mathfrak{E} \text{ then SOME } \mathfrak{F}. \mathfrak{E} = \text{kf-apply } \mathfrak{F} \text{ else } 0) \rangle$

lemma *kf-apply-km-some-kraus-family*[*simp*]:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$
shows $\langle \text{kf-apply } (\text{km-some-kraus-family } \mathfrak{E}) = \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-some-kraus-family-invalid*:
assumes $\langle \neg \text{kraus-map } \mathfrak{E} \rangle$
shows $\langle \text{km-some-kraus-family } \mathfrak{E} = 0 \rangle$
 $\langle \text{proof} \rangle$

definition *km-operators-in* :: $\langle ((\text{'a}::\text{chilbert-space}, \text{'a}) \text{ trace-class} \Rightarrow (\text{'b}::\text{chilbert-space}, \text{'b}) \text{ trace-class}) \Rightarrow (\text{'a} \Rightarrow_{CL} \text{'b}) \text{ set} \Rightarrow \text{bool} \rangle$ **where**
 $\langle \text{km-operators-in } \mathfrak{E} S \longleftrightarrow (\exists \mathfrak{F} :: (\text{'a}, \text{'b}, \text{unit}) \text{ kraus-family}. \text{kf-apply } \mathfrak{F} = \mathfrak{E} \wedge \text{kf-operators } \mathfrak{F} \subseteq S) \rangle$

lemma *km-operators-in-mono*: $\langle S \subseteq T \Longrightarrow \text{km-operators-in } \mathfrak{E} S \Longrightarrow \text{km-operators-in } \mathfrak{E} T \rangle$
 $\langle \text{proof} \rangle$

lemma *km-operators-in-kf-apply*:
assumes $\langle \text{span } (\text{kf-operators } \mathfrak{E}) \subseteq S \rangle$
shows $\langle \text{km-operators-in } (\text{kf-apply } \mathfrak{E}) S \rangle$
 $\langle \text{proof} \rangle$

lemma *km-operators-in-kf-apply-flattened*:
fixes $\mathfrak{E} :: \langle (\text{'a}::\text{chilbert-space}, \text{'b}::\text{chilbert-space}, \text{'x}::\text{CARD-1}) \text{ kraus-family} \rangle$
assumes $\langle \text{kf-operators } \mathfrak{E} \subseteq S \rangle$
shows $\langle \text{km-operators-in } (\text{kf-apply } \mathfrak{E}) S \rangle$
 $\langle \text{proof} \rangle$

lemma *km-commute*:
assumes $\langle \text{km-operators-in } \mathfrak{E} S \rangle$
assumes $\langle \text{km-operators-in } \mathfrak{F} T \rangle$
assumes $\langle S \subseteq \text{commutant } T \rangle$
shows $\langle \mathfrak{F} \circ \mathfrak{E} = \mathfrak{E} \circ \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-operators-in-UNIV*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$
shows $\langle \text{km-operators-in } \mathfrak{E} \text{ UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *separating-kraus-map-bounded-clinear*:
fixes $S :: \langle ('a::\text{hilbert-space}, 'a) \text{ trace-class set} \rangle$
assumes $\langle \text{separating-set } (\text{bounded-clinear} :: (- \Rightarrow ('b::\text{hilbert-space}, 'b) \text{ trace-class}) \Rightarrow -) S \rangle$
shows $\langle \text{separating-set } (\text{kraus-map} :: (- \Rightarrow ('b::\text{hilbert-space}, 'b) \text{ trace-class}) \Rightarrow -) S \rangle$
 $\langle \text{proof} \rangle$

3.2 Bound and norm

definition $\text{km-bound} :: \langle (('a::\text{hilbert-space}, 'a) \text{ trace-class} \Rightarrow ('b::\text{hilbert-space}, 'b) \text{ trace-class}) \Rightarrow ('a, 'a) \text{ cblinfun} \rangle$ **where**
 $\langle \text{km-bound } \mathfrak{E} = (\text{if } \exists \mathfrak{E}' :: (-, -, \text{unit}) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{E}' \text{ then kf-bound } (\text{SOME } \mathfrak{E}' :: (-, -, \text{unit}) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{E}') \text{ else } 0) \rangle$

lemma *km-bound-kf-bound*:
assumes $\langle \mathfrak{E} = \text{kf-apply } \mathfrak{F} \rangle$
shows $\langle \text{km-bound } \mathfrak{E} = \text{kf-bound } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

definition $\text{km-norm} :: \langle (('a::\text{hilbert-space}, 'a) \text{ trace-class} \Rightarrow ('b::\text{hilbert-space}, 'b) \text{ trace-class}) \Rightarrow \text{real} \rangle$ **where**
 $\langle \text{km-norm } \mathfrak{E} = \text{norm } (\text{km-bound } \mathfrak{E}) \rangle$

lemma *km-norm-kf-norm*:
assumes $\langle \mathfrak{E} = \text{kf-apply } \mathfrak{F} \rangle$
shows $\langle \text{km-norm } \mathfrak{E} = \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-invalid*:
assumes $\langle \neg \text{kraus-map } \mathfrak{E} \rangle$
shows $\langle \text{km-bound } \mathfrak{E} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-invalid*:
assumes $\langle \neg \text{kraus-map } \mathfrak{E} \rangle$
shows $\langle \text{km-norm } \mathfrak{E} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-geq0[iff]*: $\langle \text{km-norm } \mathfrak{E} \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-pos[iff]*: $\langle \text{km-bound } \mathfrak{E} \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bounded-pos*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \varrho \geq 0 \rangle$

shows $\langle \text{norm } (\mathfrak{E} \varrho) \leq \text{km-norm } \mathfrak{E} * \text{norm } \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *km-bounded*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$

shows $\langle \text{norm } (\mathfrak{E} \varrho) \leq 4 * \text{km-norm } \mathfrak{E} * \text{norm } \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-from-map*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$

shows $\langle \psi \cdot_C \text{km-bound } \mathfrak{E} \varphi = \text{trace-tc } (\mathfrak{E} (\text{tc-butterfly } \varphi \psi)) \rangle$

$\langle \text{proof} \rangle$

lemma *trace-from-km-bound*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$

shows $\langle \text{trace-tc } (\mathfrak{E} \varrho) = \text{trace-tc } (\text{compose-tcr } (\text{km-bound } \mathfrak{E}) \varrho) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-selfadjoint[iff]*: $\langle \text{selfadjoint } (\text{km-bound } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-leq-km-norm-id*: $\langle \text{km-bound } \mathfrak{E} \leq \text{km-norm } \mathfrak{E} *_R \text{id-cblinfun} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-km-some-kraus-family[simp]*: $\langle \text{kf-norm } (\text{km-some-kraus-family } \mathfrak{E}) = \text{km-norm } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

3.3 Basic Kraus maps

Zero map and constant maps. Addition and rescaling and composition of maps.

lemma *kraus-map-0[iff]*: $\langle \text{kraus-map } 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-0'[iff]*: $\langle \text{kraus-map } (\lambda -. 0) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-0[simp]*: $\langle \text{km-bound } 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *km-norm-0[simp]*: $\langle \text{km-norm } 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *km-some-kraus-family-0[simp]*: $\langle \text{km-some-kraus-family } 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-id*[*iff*]: $\langle \text{kraus-map id} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-id*[*simp*]: $\langle \text{km-bound id} = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-id-leq1*[*iff*]: $\langle \text{km-norm id} \leq 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-id-eq1*[*simp*]: $\langle \text{km-norm (id :: ('a :: \{chilbert-space, not-singleton\}, 'a) \text{ trace-class} \Rightarrow -) = 1} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-operators-in-id*[*iff*]: $\langle \text{km-operators-in id } \{ \text{id-cblinfun} \} \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-add*[*iff*]:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{kraus-map } (\lambda \varrho. \mathfrak{E} \varrho + \mathfrak{F} \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-plus'*[*iff*]:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{kraus-map } (\mathfrak{E} + \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-plus*:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-bound } (\mathfrak{E} + \mathfrak{F}) = \text{km-bound } \mathfrak{E} + \text{km-bound } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-triangle*:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-norm } (\mathfrak{E} + \mathfrak{F}) \leq \text{km-norm } \mathfrak{E} + \text{km-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-constant*[*iff*]: $\langle \text{kraus-map } (\lambda \sigma. \text{trace-tc } \sigma *_C \varrho) \rangle$ **if** $\langle \varrho \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-constant-invalid*:
 $\langle \neg \text{kraus-map } (\lambda \sigma :: ('a :: \{chilbert-space, not-singleton\}, 'a) \text{ trace-class. trace-tc } \sigma *_C \varrho) \rangle$ **if** $\langle \sim \varrho \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-scale*:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle c \geq 0 \rangle$
shows $\langle \text{kraus-map } (\lambda \varrho. c *_R \mathfrak{E} \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-scale[simp]*: $\langle km_bound (\lambda \varrho. c *_R \mathfrak{E} \varrho) = c *_R km_bound \mathfrak{E} \rangle$ **if** $\langle c \geq 0 \rangle$
 $\langle proof \rangle$

lemma *km-norm-scale[simp]*: $\langle km_norm (\lambda \varrho. c *_R \mathfrak{E} \varrho) = c * km_norm \mathfrak{E} \rangle$ **if** $\langle c \geq 0 \rangle$
 $\langle proof \rangle$

lemma *kraus-map-sandwich[iff]*: $\langle kraus_map (sandwich_tc A) \rangle$
 $\langle proof \rangle$

lemma *km-bound-sandwich[simp]*: $\langle km_bound (sandwich_tc A) = A *_R o_{CL} A \rangle$
 $\langle proof \rangle$

lemma *km-norm-sandwich[simp]*: $\langle km_norm (sandwich_tc A) = (norm A)^2 \rangle$
 $\langle proof \rangle$

lemma *km-operators-in-sandwich*: $\langle km_operators_in (sandwich_tc U) \{U\} \rangle$
 $\langle proof \rangle$

lemma *km-constant-bound[simp]*: $\langle km_bound (\lambda \sigma. trace_tc \sigma *_C \varrho) = norm \varrho *_R id_cblinfun \rangle$ **if** $\langle \varrho \geq 0 \rangle$
 $\langle proof \rangle$

lemma *km-constant-norm[simp]*: $\langle km_norm (\lambda \sigma :: ('a :: \{chilbert_space, not_singleton\}, 'a) trace_class. trace_tc \sigma *_C \varrho) = norm \varrho \rangle$ **if** $\langle \varrho \geq 0 \rangle$
 $\langle proof \rangle$

lemma *km-constant-norm-leq[simp]*: $\langle km_norm (\lambda \sigma :: ('a :: chilbert_space, 'a) trace_class. trace_tc \sigma *_C \varrho) \leq norm \varrho \rangle$
 $\langle proof \rangle$

lemma *kraus-map-comp*:
assumes $\langle kraus_map \mathfrak{E} \rangle$ **and** $\langle kraus_map \mathfrak{F} \rangle$
shows $\langle kraus_map (\mathfrak{E} \circ \mathfrak{F}) \rangle$
 $\langle proof \rangle$

lemma *km-comp-norm-leq*:
assumes $\langle kraus_map \mathfrak{E} \rangle$ **and** $\langle kraus_map \mathfrak{F} \rangle$
shows $\langle km_norm (\mathfrak{E} \circ \mathfrak{F}) \leq km_norm \mathfrak{E} * km_norm \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *km-bound-comp-sandwich*:
assumes $\langle kraus_map \mathfrak{E} \rangle$
shows $\langle km_bound (\lambda \varrho. \mathfrak{E} (sandwich_tc U \varrho)) = sandwich (U*) (km_bound \mathfrak{E}) \rangle$
 $\langle proof \rangle$

lemma *km-norm-comp-sandwich-coiso*:
assumes $\langle \text{isometry } (U*) \rangle$
shows $\langle \text{km-norm } (\lambda \varrho. \mathfrak{E} (\text{sandwich-tc } U \ \varrho)) = \text{km-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-comp-sandwich-iso*:
assumes $\langle \text{isometry } U \rangle$
shows $\langle \text{km-bound } (\lambda \varrho. \text{sandwich-tc } U \ (\mathfrak{E} \ \varrho)) = \text{km-bound } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-comp-sandwich-iso*:
assumes $\langle \text{isometry } U \rangle$
shows $\langle \text{km-norm } (\lambda \varrho. \text{sandwich-tc } U \ (\mathfrak{E} \ \varrho)) = \text{km-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-sum*:
assumes $\langle \bigwedge x. x \in A \implies \text{kraus-map } (\mathfrak{E} \ x) \rangle$
shows $\langle \text{kraus-map } (\sum x \in A. \mathfrak{E} \ x) \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-sum*:
assumes $\langle \bigwedge x. x \in A \implies \text{kraus-map } (\mathfrak{E} \ x) \rangle$
shows $\langle \text{km-bound } (\sum x \in A. \mathfrak{E} \ x) = (\sum x \in A. \text{km-bound } (\mathfrak{E} \ x)) \rangle$
 $\langle \text{proof} \rangle$

3.4 Infinite sums

lemma
assumes $\langle \bigwedge \varrho. ((\lambda a. \text{sandwich-tc } (f \ a) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ A \rangle$
defines $\langle EE \equiv \text{Set.filter } (\lambda (E, -). E \neq 0) ((\lambda a. (f \ a, \ a)) \text{ ` } A) \rangle$
shows *kraus-mapI-sum*: $\langle \text{kraus-map } \mathfrak{E} \rangle$
and *kraus-map-sum-kraus-family*: $\langle \text{kraus-family } EE \rangle$
and *kraus-map-sum-kf-apply*: $\langle \mathfrak{E} = \text{kf-apply } (\text{Abs-kraus-family } EE) \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-infsum-sandwich*:
assumes $\langle \bigwedge \varrho. (\lambda a. \text{sandwich-tc } (f \ a) \ \varrho) \text{ summable-on } A \rangle$
shows $\langle \text{kraus-map } (\lambda \varrho. \sum_{\infty a \in A. \text{sandwich-tc } (f \ a) \ \varrho}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-sum-sandwich*: $\langle \text{kraus-map } (\lambda \varrho. \sum a \in A. \text{sandwich-tc } (f \ a) \ \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-as-infsum*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$

shows $\langle \exists M. \forall \varrho. ((\lambda E. \text{sandwich-tc } E \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ M \rangle$

$\langle \text{proof} \rangle$

definition *km-summable* :: $\langle ('a \Rightarrow ('b::\text{chilbert-space}, 'b) \text{ trace-class} \Rightarrow ('c::\text{chilbert-space}, 'c) \text{ trace-class}) \Rightarrow 'a \text{ set} \Rightarrow \text{bool} \rangle$ **where**

$\langle \text{km-summable } \mathfrak{E} \ A \longleftrightarrow \text{summable-on-in cweak-operator-topology } (\lambda a. \text{km-bound } (\mathfrak{E} \ a)) \ A \rangle$

lemma *km-summable-kf-summable*:

assumes $\langle \bigwedge a \ \varrho. a \in A \implies \mathfrak{E} \ a \ \varrho = \mathfrak{F} \ a \ *_{kr} \ \varrho \rangle$

shows $\langle \text{km-summable } \mathfrak{E} \ A \longleftrightarrow \text{kf-summable } \mathfrak{F} \ A \rangle$

$\langle \text{proof} \rangle$

lemma *km-summable-summable*:

assumes *km*: $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} \ a) \rangle$

assumes *sum*: $\langle \text{km-summable } \mathfrak{E} \ A \rangle$

shows $\langle (\lambda a. \mathfrak{E} \ a \ \varrho) \text{ summable-on } A \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-infsum*:

assumes *km*: $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} \ a) \rangle$

assumes *sum*: $\langle \text{km-summable } \mathfrak{E} \ A \rangle$

shows $\langle \text{kraus-map } (\lambda \varrho. \sum_{\infty} a \in A. \mathfrak{E} \ a \ \varrho) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-infsum*:

assumes *km*: $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} \ a) \rangle$

assumes *sum*: $\langle \text{km-summable } \mathfrak{E} \ A \rangle$

shows $\langle \text{km-bound } (\lambda \varrho. \sum_{\infty} a \in A. \mathfrak{E} \ a \ \varrho) = \text{infsum-in cweak-operator-topology } (\lambda a. \text{km-bound } (\mathfrak{E} \ a)) \ A \rangle$

$\langle \text{proof} \rangle$

lemma *km-norm-infsum*:

assumes *km*: $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} \ a) \rangle$

assumes *sum*: $\langle (\lambda a. \text{km-norm } (\mathfrak{E} \ a)) \text{ summable-on } A \rangle$

shows $\langle \text{km-norm } (\lambda \varrho. \sum_{\infty} a \in A. \mathfrak{E} \ a \ \varrho) \leq (\sum_{\infty} a \in A. \text{km-norm } (\mathfrak{E} \ a)) \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-has-sum*:

assumes $\langle \bigwedge x. x \in A \implies \text{kraus-map } (\mathfrak{E} \ x) \rangle$

assumes $\langle \text{km-summable } \mathfrak{E} \ A \rangle$

assumes $\langle (\mathfrak{E} \text{ has-sum } \mathfrak{F}) \ A \rangle$

shows $\langle \text{kraus-map } \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

lemma *km-summable-iff-sums-to-kraus-map*:

assumes $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} a) \rangle$

shows $\langle \text{km-summable } \mathfrak{E} A \longleftrightarrow (\exists \mathfrak{F}. (\forall t. ((\lambda x. \mathfrak{E} x t) \text{ has-sum } \mathfrak{F} t) A) \wedge \text{kraus-map } \mathfrak{F}) \rangle$

$\langle \text{proof} \rangle$

3.5 Tensor products

definition *km-tensor-exists* :: $\langle ((\text{'a ell2}, \text{'b ell2}) \text{ trace-class} \Rightarrow (\text{'c ell2}, \text{'d ell2}) \text{ trace-class})$
 $\Rightarrow ((\text{'e ell2}, \text{'f ell2}) \text{ trace-class} \Rightarrow (\text{'g ell2}, \text{'h ell2}) \text{ trace-class}) \Rightarrow \text{bool} \rangle$

where

$\langle \text{km-tensor-exists } \mathfrak{E} \mathfrak{F} \longleftrightarrow (\exists \mathfrak{E}\mathfrak{F}. \text{ bounded-clinear } \mathfrak{E}\mathfrak{F} \wedge (\forall \varrho \sigma. \mathfrak{E}\mathfrak{F} (\text{tc-tensor } \varrho \sigma) = \text{tc-tensor } (\mathfrak{E} \varrho) (\mathfrak{F} \sigma))) \rangle$

definition *km-tensor* :: $\langle ((\text{'a ell2}, \text{'c ell2}) \text{ trace-class} \Rightarrow (\text{'e ell2}, \text{'g ell2}) \text{ trace-class})$
 $\Rightarrow ((\text{'b ell2}, \text{'d ell2}) \text{ trace-class} \Rightarrow (\text{'f ell2}, \text{'h ell2}) \text{ trace-class})$
 $\Rightarrow ((\text{'a} \times \text{'b}) \text{ ell2}, (\text{'c} \times \text{'d}) \text{ ell2}) \text{ trace-class} \Rightarrow ((\text{'e} \times \text{'f}) \text{ ell2}, (\text{'g} \times \text{'h}) \text{ ell2})$

trace-class **where**

$\langle \text{km-tensor } \mathfrak{E} \mathfrak{F} = (\text{if km-tensor-exists } \mathfrak{E} \mathfrak{F}$
 $\text{ then SOME } \mathfrak{E}\mathfrak{F}. \text{ bounded-clinear } \mathfrak{E}\mathfrak{F} \wedge (\forall \varrho \sigma. \mathfrak{E}\mathfrak{F} (\text{tc-tensor } \varrho \sigma) = \text{tc-tensor } (\mathfrak{E} \varrho) (\mathfrak{F} \sigma))$
 $\text{ else } 0) \rangle$

lemma *km-tensor-invalid*:

assumes $\langle \neg \text{km-tensor-exists } \mathfrak{E} \mathfrak{F} \rangle$

shows $\langle \text{km-tensor } \mathfrak{E} \mathfrak{F} = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-exists-bounded-clinear*[*iff*]:

assumes $\langle \text{km-tensor-exists } \mathfrak{E} \mathfrak{F} \rangle$

shows $\langle \text{bounded-clinear } (\text{km-tensor } \mathfrak{E} \mathfrak{F}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-apply*[*simp*]:

assumes $\langle \text{km-tensor-exists } \mathfrak{E} \mathfrak{F} \rangle$

shows $\langle \text{km-tensor } \mathfrak{E} \mathfrak{F} (\text{tc-tensor } \varrho \sigma) = \text{tc-tensor } (\mathfrak{E} \varrho) (\mathfrak{F} \sigma) \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-unique*:

assumes $\langle \text{bounded-clinear } \mathfrak{E}\mathfrak{F} \rangle$

assumes $\langle \bigwedge \varrho \sigma. \mathfrak{E}\mathfrak{F} (\text{tc-tensor } \varrho \sigma) = \text{tc-tensor } (\mathfrak{E} \varrho) (\mathfrak{F} \sigma) \rangle$

shows $\langle \mathfrak{E}\mathfrak{F} = \text{km-tensor } \mathfrak{E} \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-kf-tensor*: $\langle \text{km-tensor } (\text{kf-apply } \mathfrak{E}) (\text{kf-apply } \mathfrak{F}) = \text{kf-apply } (\text{kf-tensor } \mathfrak{E} \mathfrak{F}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-kraus-map*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{kraus-map } (\text{km-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-kraus-map-exists*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-tensor-exists } \mathfrak{E} \ \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-as-infsum*:

assumes $\langle \bigwedge \varrho. ((\lambda i. \text{sandwich-tc } (E \ i) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ I \rangle$
assumes $\langle \bigwedge \varrho. ((\lambda j. \text{sandwich-tc } (F \ j) \ \varrho) \text{ has-sum } \mathfrak{F} \ \varrho) \ J \rangle$
shows $\langle \text{km-tensor } \mathfrak{E} \ \mathfrak{F} \ \varrho = (\sum_{\infty (i,j) \in I \times J. \text{sandwich-tc } (E \ i \otimes_o F \ j) \ \varrho} \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-tensor*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-bound } (\text{km-tensor } \mathfrak{E} \ \mathfrak{F}) = \text{km-bound } \mathfrak{E} \otimes_o \text{km-bound } \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

lemma *km-norm-tensor*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-norm } (\text{km-tensor } \mathfrak{E} \ \mathfrak{F}) = \text{km-norm } \mathfrak{E} * \text{km-norm } \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-compose-distrib*:

assumes $\langle \text{km-tensor-exists } \mathfrak{E} \ \mathfrak{G} \rangle$ **and** $\langle \text{km-tensor-exists } \mathfrak{F} \ \mathfrak{H} \rangle$
shows $\langle \text{km-tensor } (\mathfrak{E} \ o \ \mathfrak{F}) \ (\mathfrak{G} \ o \ \mathfrak{H}) = \text{km-tensor } \mathfrak{E} \ \mathfrak{G} \ o \ \text{km-tensor } \mathfrak{F} \ \mathfrak{H} \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-tensor-right[simp]*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kraus-map } (\lambda \sigma. \text{tc-tensor } \sigma \ \varrho) \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-tensor-left[simp]*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kraus-map } (\lambda \sigma. \text{tc-tensor } \varrho \ \sigma) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-tensor-right[simp]*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{km-bound } (\lambda \sigma. \text{tc-tensor } \sigma \ \varrho) = \text{norm } \varrho *_C \text{id-cblinfun} \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-tensor-left[simp]*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{km-bound } (\lambda \sigma. \text{tc-tensor } \varrho \ \sigma) = \text{norm } \varrho *_C \text{id-cblinfun} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-tensor-right[simp]*:

assumes $\langle \varrho \geq 0 \rangle$

shows $\langle \text{km-norm } (\lambda \sigma. \text{tc-tensor } \sigma \ \varrho) = \text{norm } \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-tensor-left[simp]*:

assumes $\langle \varrho \geq 0 \rangle$

shows $\langle \text{km-norm } (\lambda \sigma. \text{tc-tensor } \varrho \ \sigma) = \text{norm } \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *km-operators-in-tensor*:

assumes $\langle \text{km-operators-in } \mathfrak{E} \ S \rangle$

assumes $\langle \text{km-operators-in } \mathfrak{F} \ T \rangle$

shows $\langle \text{km-operators-in } (\text{km-tensor } \mathfrak{E} \ \mathfrak{F}) \ (\text{span } \{s \otimes_o t \mid s \in S \wedge t \in T\}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-sandwich-tc*:

$\langle \text{km-tensor } (\text{sandwich-tc } A) \ (\text{sandwich-tc } B) = \text{sandwich-tc } (A \otimes_o B) \rangle$

$\langle \text{proof} \rangle$

3.6 Trace and partial trace

definition $\langle \text{km-trace-preserving } \mathfrak{E} \longleftrightarrow (\exists \mathfrak{F}::(-, -, \text{unit}) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \wedge \text{kf-trace-preserving } \mathfrak{F}) \rangle$

lemma *km-trace-preserving-def'*: $\langle \text{km-trace-preserving } \mathfrak{E} \longleftrightarrow (\exists \mathfrak{F}::(-, -, 'c) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \wedge \text{kf-trace-preserving } \mathfrak{F}) \rangle$

— Has a more general type than *km-trace-preserving-def*

$\langle \text{proof} \rangle$

definition *km-trace-reducing-def*: $\langle \text{km-trace-reducing } \mathfrak{E} \longleftrightarrow (\exists \mathfrak{F}::(-, -, \text{unit}) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \wedge \text{kf-trace-reducing } \mathfrak{F}) \rangle$

lemma *km-trace-reducing-def'*: $\langle \text{km-trace-reducing } \mathfrak{E} \longleftrightarrow (\exists \mathfrak{F}::(-, -, 'c) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \wedge \text{kf-trace-reducing } \mathfrak{F}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-trace-preserving-apply[simp]*: $\langle \text{km-trace-preserving } (\text{kf-apply } \mathfrak{E}) = \text{kf-trace-preserving } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *km-trace-reducing-apply[simp]*: $\langle \text{km-trace-reducing } (\text{kf-apply } \mathfrak{E}) = \text{kf-trace-reducing } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *km-trace-preserving-iff*: $\langle \text{km-trace-preserving } \mathfrak{E} \longleftrightarrow \text{kraus-map } \mathfrak{E} \wedge (\forall \varrho. \text{trace-tc } (\mathfrak{E} \ \varrho) = \text{trace-tc } \varrho) \rangle$

$\langle \text{proof} \rangle$

lemma *km-trace-reducing-iff*: $\langle km\text{-}trace\text{-}reducing\ \mathfrak{E} \longleftrightarrow kraus\text{-}map\ \mathfrak{E} \wedge (\forall \varrho \geq 0. trace\text{-}tc\ (\mathfrak{E}\ \varrho) \leq trace\text{-}tc\ \varrho) \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-imp-reducing*:
assumes $\langle km\text{-}trace\text{-}preserving\ \mathfrak{E} \rangle$
shows $\langle km\text{-}trace\text{-}reducing\ \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-id[iff]*: $\langle km\text{-}trace\text{-}preserving\ id \rangle$
 $\langle proof \rangle$

lemma *km-trace-reducing-iff-norm-leq1*: $\langle km\text{-}trace\text{-}reducing\ \mathfrak{E} \longleftrightarrow kraus\text{-}map\ \mathfrak{E} \wedge km\text{-}norm\ \mathfrak{E} \leq 1 \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-iff-bound-id*: $\langle km\text{-}trace\text{-}preserving\ \mathfrak{E} \longleftrightarrow kraus\text{-}map\ \mathfrak{E} \wedge km\text{-}bound\ \mathfrak{E} = id\text{-}cblinfun \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-iff-bound-id'*:
fixes $\mathfrak{E} :: \langle ('a :: \{chilbert\text{-}space, not\text{-}singleton\}, 'a)\ trace\text{-}class \Rightarrow - \rangle$
shows $\langle km\text{-}trace\text{-}preserving\ \mathfrak{E} \longleftrightarrow km\text{-}bound\ \mathfrak{E} = id\text{-}cblinfun \rangle$
 $\langle proof \rangle$

lemma *km-trace-norm-preserving*: $\langle km\text{-}norm\ \mathfrak{E} \leq 1 \rangle$ **if** $\langle km\text{-}trace\text{-}preserving\ \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *km-trace-norm-preserving-eq*:
fixes $\mathfrak{E} :: \langle ('a :: \{chilbert\text{-}space, not\text{-}singleton\}, 'a)\ trace\text{-}class \Rightarrow ('b :: chilbert\text{-}space, 'b)\ trace\text{-}class) \rangle$
assumes $\langle km\text{-}trace\text{-}preserving\ \mathfrak{E} \rangle$
shows $\langle km\text{-}norm\ \mathfrak{E} = 1 \rangle$
 $\langle proof \rangle$

lemma *kraus-map-trace*: $\langle kraus\text{-}map\ (one\text{-}dim\text{-}iso\ o\ trace\text{-}tc) \rangle$
 $\langle proof \rangle$

lemma *trace-preserving-trace-kraus-map[iff]*: $\langle km\text{-}trace\text{-}preserving\ (one\text{-}dim\text{-}iso\ o\ trace\text{-}tc) \rangle$
 $\langle proof \rangle$

lemma *km-trace-bound[simp]*: $\langle km\text{-}bound\ (one\text{-}dim\text{-}iso\ o\ trace\text{-}tc) = id\text{-}cblinfun \rangle$
 $\langle proof \rangle$

lemma *km-trace-norm-eq1[simp]*: $\langle km\text{-}norm\ (one\text{-}dim\text{-}iso\ o\ trace\text{-}tc :: ('a :: \{chilbert\text{-}space, not\text{-}singleton\}, 'a)\ trace\text{-}class \Rightarrow -) = 1 \rangle$
 $\langle proof \rangle$

lemma *km-trace-norm-leq1*[simp]: $\langle km\text{-}norm\ (one\text{-}dim\text{-}iso\ o\ trace\text{-}tc) \leq 1 \rangle$
 $\langle proof \rangle$

lemma *kraus-map-partial-trace*[iff]: $\langle kraus\text{-}map\ partial\text{-}trace \rangle$
 $\langle proof \rangle$

lemma *partial-trace-ignores-kraus-map*:
assumes $\langle km\text{-}trace\text{-}preserving\ \mathfrak{E} \rangle$
assumes $\langle kraus\text{-}map\ \mathfrak{F} \rangle$
shows $\langle partial\text{-}trace\ (km\text{-}tensor\ \mathfrak{F}\ \mathfrak{E}\ \varrho) = \mathfrak{F}\ (partial\text{-}trace\ \varrho) \rangle$
 $\langle proof \rangle$

lemma *km-partial-trace-bound*[simp]: $\langle km\text{-}bound\ partial\text{-}trace = id\text{-}cblinfun \rangle$
 $\langle proof \rangle$

lemma *km-partial-trace-norm*[simp]:
shows $\langle km\text{-}norm\ partial\text{-}trace = 1 \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-tensor*:
assumes $\langle km\text{-}trace\text{-}preserving\ \mathfrak{E} \rangle$ **and** $\langle km\text{-}trace\text{-}preserving\ \mathfrak{F} \rangle$
shows $\langle km\text{-}trace\text{-}preserving\ (km\text{-}tensor\ \mathfrak{E}\ \mathfrak{F}) \rangle$
 $\langle proof \rangle$

lemma *km-trace-reducing-tensor*:
assumes $\langle km\text{-}trace\text{-}reducing\ \mathfrak{E} \rangle$ **and** $\langle km\text{-}trace\text{-}reducing\ \mathfrak{F} \rangle$
shows $\langle km\text{-}trace\text{-}reducing\ (km\text{-}tensor\ \mathfrak{E}\ \mathfrak{F}) \rangle$
 $\langle proof \rangle$

3.7 Complete measurements

definition $\langle km\text{-}complete\text{-}measurement\ B\ \varrho = (\sum_{x \in B} sandwich\text{-}tc\ (selfbutter\ (sgn\ x))\ \varrho) \rangle$
abbreviation $\langle km\text{-}complete\text{-}measurement\text{-}ket \equiv km\text{-}complete\text{-}measurement\ (range\ ket) \rangle$

lemma *km-complete-measurement-kf-complete-measurement*: $\langle km\text{-}complete\text{-}measurement\ B = kf\text{-}apply\ (kf\text{-}complete\text{-}measurement\ B) \rangle$ **if** $\langle is\text{-}ortho\text{-}set\ B \rangle$
 $\langle proof \rangle$

lemma *km-complete-measurement-ket-kf-complete-measurement-ket*: $\langle km\text{-}complete\text{-}measurement\text{-}ket = kf\text{-}apply\ kf\text{-}complete\text{-}measurement\text{-}ket \rangle$
 $\langle proof \rangle$

lemma *km-complete-measurement-has-sum*:
assumes $\langle is\text{-}ortho\text{-}set\ B \rangle$
shows $\langle ((\lambda x. sandwich\text{-}tc\ (selfbutter\ (sgn\ x))\ \varrho)\ has\text{-}sum\ km\text{-}complete\text{-}measurement\ B\ \varrho)\ B \rangle$

⟨proof⟩

lemma *km-complete-measurement-ket-has-sum*:

⟨((λx. sandwich-tc (selfbutter (ket x)) ρ) has-sum km-complete-measurement-ket ρ) UNIV⟩

⟨proof⟩

lemma *km-bound-complete-measurement*:

assumes ⟨is-ortho-set B⟩

shows ⟨km-bound (km-complete-measurement B) ≤ id-cblinfun⟩

⟨proof⟩

lemma *km-norm-complete-measurement*:

assumes ⟨is-ortho-set B⟩

shows ⟨km-norm (km-complete-measurement B) ≤ 1⟩

⟨proof⟩

lemma *km-bound-complete-measurement-onb[simp]*:

assumes ⟨is-onb B⟩

shows ⟨km-bound (km-complete-measurement B) = id-cblinfun⟩

⟨proof⟩

lemma *km-bound-complete-measurement-ket[simp]*: ⟨km-bound km-complete-measurement-ket = id-cblinfun⟩

⟨proof⟩

lemma *km-norm-complete-measurement-onb[simp]*:

fixes B :: ⟨'a::{not-singleton, hilbert-space} set⟩

assumes ⟨is-onb B⟩

shows ⟨km-norm (km-complete-measurement B) = 1⟩

⟨proof⟩

lemma *km-norm-complete-measurement-ket[simp]*:

shows ⟨km-norm km-complete-measurement-ket = 1⟩

⟨proof⟩

lemma *kraus-map-complete-measurement*:

assumes ⟨is-ortho-set B⟩

shows ⟨kraus-map (km-complete-measurement B)⟩

⟨proof⟩

lemma *kraus-map-complete-measurement-ket[iff]*:

shows ⟨kraus-map km-complete-measurement-ket⟩

⟨proof⟩

lemma *km-complete-measurement-idem[simp]*:

assumes ⟨is-ortho-set B⟩

shows ⟨km-complete-measurement B (km-complete-measurement B ρ) = km-complete-measurement B ρ⟩

⟨proof⟩

lemma *km-complete-measurement-ket-idem*[simp]:
 $\langle km_complete_measurement_ket (km_complete_measurement_ket \varrho) = km_complete_measurement_ket \varrho \rangle$
 ⟨proof⟩

lemma *km-complete-measurement-has-sum-onb*:
assumes ⟨is-onb B⟩
shows ⟨((λx. sandwich-tc (selfbutter x) ρ) has-sum km-complete-measurement B ρ) B⟩
 ⟨proof⟩

lemma *km-complete-measurement-ket-diagonal-operator*[simp]:
 $\langle km_complete_measurement_ket (diagonal_operator_tc f) = diagonal_operator_tc f \rangle$
 ⟨proof⟩

lemma *km-operators-complete-measurement*:
assumes ⟨is-ortho-set B⟩
shows ⟨km-operators-in (km-complete-measurement B) (span (selfbutter ‘ B))⟩
 ⟨proof⟩

lemma *km-operators-complete-measurement-ket*:
shows ⟨km-operators-in km-complete-measurement-ket (span (range (λc. (selfbutter (ket c))))))⟩
 ⟨proof⟩

lemma *km-complete-measurement-ket-butterket*[simp]:
 $\langle km_complete_measurement_ket (tc_butterfly (ket c) (ket c)) = tc_butterfly (ket c) (ket c) \rangle$
 ⟨proof⟩

lemma *km-complete-measurement-tensor*:
assumes ⟨is-ortho-set B⟩ **and** ⟨is-ortho-set C⟩
shows ⟨km-tensor (km-complete-measurement B) (km-complete-measurement C)
 $= km_complete_measurement ((\lambda(b,c). b \otimes_s c) ‘ (B \times C)) \rangle$
 ⟨proof⟩

lemma *km-complete-measurement-ket-tensor*:
shows ⟨km-tensor (km-complete-measurement-ket :: (‘a ell2, -) trace-class ⇒ -) (km-complete-measurement-ket
 :: (‘b ell2, -) trace-class ⇒ -)
 $= km_complete_measurement_ket \rangle$
 ⟨proof⟩

lemma *km-tensor-0-left*[simp]: ⟨km-tensor (0 :: (‘a ell2, ‘b ell2) trace-class ⇒ (‘c ell2, ‘d ell2)
 trace-class) $\mathfrak{E} = 0 \rangle$
 ⟨proof⟩

lemma *km-tensor-0-right*[simp]: ⟨km-tensor $\mathfrak{E}$ (0 :: (‘a ell2, ‘b ell2) trace-class ⇒ (‘c ell2, ‘d
 ell2) trace-class) = 0⟩
 ⟨proof⟩

unbundle *no kraus-map-syntax*
unbundle *no cblinfun-syntax*

end

References

- [1] K. Kraus. *States, effects, and operations: fundamental notions of quantum theory*, volume 190 of *LNP*. Springer, 1983. Defines Kraus maps in Section 3, Theorem 1. Only considers separable Hilbert spaces.