

The Kolmogorov-Chentsov Theorem

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October 13, 2025

Abstract

Continuous-time stochastic processes often carry the condition of having almost-surely continuous paths. If some process X satisfies certain bounds on its expectation, then the Kolmogorov-Chentsov theorem lets us construct a modification of X , i.e. a process X' such that $\forall t. X_t = X'_t$ almost surely, that has Hölder continuous paths.

In this work, we mechanise the Kolmogorov-Chentsov theorem. To get there, we develop a theory of stochastic processes, together with Hölder continuity, convergence in measure, and arbitrary intervals of dyadic rationals.

With this, we pave the way towards a construction of Brownian motion. The work is based on the exposition in Achim Klenke's probability theory text [1].

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1 Supporting lemmas

theory *Kolmogorov-Chentsov-Extras*
imports *HOL-Probability.Probability*
begin

lemma *atLeastAtMost-induct*[*consumes 1, case-names base Suc*]:
assumes $x \in \{n..m\}$
and $P\ n$
and $\bigwedge k. \llbracket k \geq n; k < m; P\ k \rrbracket \implies P\ (Suc\ k)$
shows $P\ x$
 $\langle proof \rangle$

lemma *eventually-prodI'*:
assumes *eventually* $P\ F$ *eventually* $Q\ G\ \forall x\ y. P\ x \longrightarrow Q\ y \longrightarrow R\ (x,y)$
shows *eventually* $R\ (F \times_F G)$
 $\langle proof \rangle$

Analogous to $\llbracket \text{almost-everywhere } ?M\ ?P; \bigwedge N. \llbracket \bigwedge x. x \in \text{space } ?M - N \implies ?P\ x; N \in \text{null-sets } ?M \rrbracket \implies ?thesis \rrbracket \implies ?thesis$

lemma *AE-I3*:
assumes $\bigwedge x. x \in \text{space } M - N \implies P\ x\ N \in \text{null-sets } M$
shows *AE* x *in* $M. P\ x$
 $\langle proof \rangle$

Extends $\llbracket (?f \longrightarrow ?l)\ ?F; (?g \longrightarrow ?m)\ ?F \rrbracket \implies ((\lambda x. \text{dist } (?f\ x)\ (?g\ x)) \longrightarrow \text{dist } ?l\ ?m)\ ?F$

lemma *tendsto-dist-prod*:
fixes $l\ m :: 'a :: \text{metric-space}$
assumes $f: (f \longrightarrow l)\ F$
and $g: (g \longrightarrow m)\ G$
shows $((\lambda x. \text{dist } (f\ (\text{fst } x))\ (g\ (\text{snd } x))) \longrightarrow \text{dist } l\ m)\ (F \times_F G)$
 $\langle proof \rangle$

lemma *borel-measurable-at-within-metric* [*measurable*]:
fixes $f :: 'a :: \text{first-countable-topology} \Rightarrow 'b \Rightarrow 'c :: \text{metric-space}$
assumes [*measurable*]: $\bigwedge t. t \in S \implies f\ t \in \text{borel-measurable } M$
and *convergent*: $\bigwedge \omega. \omega \in \text{space } M \implies \exists l. ((\lambda t. f\ t\ \omega) \longrightarrow l)\ (\text{at } x\ \text{within } S)$
and *nontrivial*: $\text{at } x\ \text{within } S \neq \perp$
shows $(\lambda \omega. \text{Lim } (\text{at } x\ \text{within } S)\ (\lambda t. f\ t\ \omega)) \in \text{borel-measurable } M$
 $\langle proof \rangle$

lemma *Max-finite-image-ex*:
assumes *finite* $S\ S \neq \{\}$ $P\ (\text{MAX } k \in S. f\ k)$
shows $\exists k \in S. P\ (f\ k)$
 $\langle proof \rangle$

lemma *measure-leq-emeasure-ennreal*: $0 \leq x \implies \text{emeasure } M\ A \leq \text{ennreal } x \implies \text{measure } M\ A \leq x$

<proof>

lemma *Union-add-subset*: $(m :: \text{nat}) \leq n \implies (\bigcup k. A (k + n)) \subseteq (\bigcup k. A (k + m))$
<proof>

lemma *floor-in-Nats* [simp]: $x \geq 0 \implies \lfloor x \rfloor \in \mathbb{N}$
<proof>

lemma *triangle-ineq-list*:
fixes $l :: ('a :: \text{metric-space}) \text{ list}$
assumes $l \neq []$
shows $\text{dist } (\text{hd } l) (\text{last } l) \leq (\sum_{i=1..length\ l - 1} \text{dist } (l[i]) (l[i+1]))$
<proof>

lemma *triangle-ineq-sum*:
fixes $f :: \text{nat} \Rightarrow 'a :: \text{metric-space}$
assumes $n \leq m$
shows $\text{dist } (f\ n) (f\ m) \leq (\sum_{i=\text{Suc } n..m} \text{dist } (f\ i) (f\ (i+1)))$
<proof>

lemma (in *product-prob-space*) *indep-vars-PiM-coordinate*:
assumes $I \neq \{\}$
shows $\text{prob-space.indep-vars } (\Pi_M\ i \in I. M\ i)\ M\ (\lambda x\ f. f\ x)\ I$
<proof>

lemma (in *prob-space*) *indep-sets-indep-set*:
assumes $\text{indep-sets } F\ I\ i \in I\ j \in I\ i \neq j$
shows $\text{indep-set } (F\ i)\ (F\ j)$
<proof>

lemma (in *prob-space*) *indep-vars-indep-var*:
assumes $\text{indep-vars } M'\ X\ I\ i \in I\ j \in I\ i \neq j$
shows $\text{indep-var } (M'\ i)\ (X\ i)\ (M'\ j)\ (X\ j)$
<proof>

end

2 Intervals of dyadic rationals

theory *Dyadic-Interval*
imports *HOL-Analysis.Analysis*
begin

In this file we describe intervals of dyadic numbers $S..T$ for reals $S\ T$. We use the floor and ceiling functions to approximate the numbers with increasing accuracy.

lemma *frac-floor*: $\lfloor x \rfloor = x - \text{frac } x$

$\langle \text{proof} \rangle$

lemma *frac-ceil*: $\lceil x \rceil = x + \text{frac } (-x)$
 $\langle \text{proof} \rangle$

lemma *floor-pow2-lim*: $(\lambda n. \lfloor 2^n * T \rfloor / 2^n) \longrightarrow T$
 $\langle \text{proof} \rangle$

lemma *floor-pow2-leq*: $\lfloor 2^n * T \rfloor / 2^n \leq T$
 $\langle \text{proof} \rangle$

lemma *ceil-pow2-lim*: $(\lambda n. \lceil 2^n * T \rceil / 2^n) \longrightarrow T$
 $\langle \text{proof} \rangle$

lemma *ceil-pow2-geq*: $\lceil 2^n * T \rceil / 2^n \geq T$
 $\langle \text{proof} \rangle$

dyadic_interval_step $n \ S \ T$ is the collection of dyadic numbers in $\{S..T\}$ with denominator 2^n . As $n \rightarrow \infty$ this collection approximates $\{S..T\}$. Compare with *dyadics* $\equiv \bigcup_{k \ m} \{ \text{of-nat } m / (2::?'a)^k \}$

definition *dyadic-interval-step* :: $\text{nat} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real set}$
where *dyadic-interval-step* $n \ S \ T \equiv (\lambda k. k / (2^n)) \text{ ' } \{ \lceil 2^n * S \rceil .. \lfloor 2^n * T \rfloor \}$

definition *dyadic-interval* :: $\text{real} \Rightarrow \text{real} \Rightarrow \text{real set}$
where *dyadic-interval* $S \ T \equiv (\bigcup n. \text{dyadic-interval-step } n \ S \ T)$

lemma *dyadic-interval-step-empty[simp]*: $T < S \implies \text{dyadic-interval-step } n \ S \ T = \{\}$
 $\langle \text{proof} \rangle$

lemma *dyadic-interval-step-singleton[simp]*: $X \in \mathbb{Z} \implies \text{dyadic-interval-step } n \ X \ X = \{X\}$
 $\langle \text{proof} \rangle$

lemma *dyadic-interval-step-zero [simp]*: $\text{dyadic-interval-step } 0 \ S \ T = \text{real-of-int ' } \{ \lceil S \rceil .. \lfloor T \rfloor \}$
 $\langle \text{proof} \rangle$

lemma *dyadic-interval-step-mem [intro]*:
assumes $x \geq 0 \ T \geq 0 \ x \leq T$
shows $\lfloor 2^n * x \rfloor / 2^n \in \text{dyadic-interval-step } n \ 0 \ T$
 $\langle \text{proof} \rangle$

lemma *dyadic-interval-step-iff*:
 $x \in \text{dyadic-interval-step } n \ S \ T \longleftrightarrow$
 $(\exists k. k \geq \lceil 2^n * S \rceil \wedge k \leq \lfloor 2^n * T \rfloor \wedge x = k / 2^n)$
 $\langle \text{proof} \rangle$

lemma *dyadic-interval-step-memI [intro]*:

assumes $\exists k::int. x = k/2^n \wedge x \geq S \wedge x \leq T$
shows $x \in \text{dyadic-interval-step } n \ S \ T$
 $\langle \text{proof} \rangle$

lemma *mem-dyadic-interval*: $x \in \text{dyadic-interval } S \ T \longleftrightarrow (\exists n. x \in \text{dyadic-interval-step } n \ S \ T)$
 $\langle \text{proof} \rangle$

lemma *mem-dyadic-intervalI*: $\exists n. x \in \text{dyadic-interval-step } n \ S \ T \implies x \in \text{dyadic-interval } S \ T$
 $\langle \text{proof} \rangle$

lemma *dyadic-step-leq*: $x \in \text{dyadic-interval-step } n \ S \ T \implies x \leq T$
 $\langle \text{proof} \rangle$

lemma *dyadics-leq*: $x \in \text{dyadic-interval } S \ T \implies x \leq T$
 $\langle \text{proof} \rangle$

lemma *dyadic-step-geq*: $x \in \text{dyadic-interval-step } n \ S \ T \implies x \geq S$
 $\langle \text{proof} \rangle$

lemma *dyadics-geq*: $x \in \text{dyadic-interval } S \ T \implies x \geq S$
 $\langle \text{proof} \rangle$

corollary *dyadic-interval-subset-interval* [simp]: $(\text{dyadic-interval } 0 \ T) \subseteq \{0..T\}$
 $\langle \text{proof} \rangle$

lemma *zero-in-dyadics*: $T \geq 0 \implies 0 \in \text{dyadic-interval-step } n \ 0 \ T$
 $\langle \text{proof} \rangle$

The following theorem is useful for reasoning with `at_within`

lemma *dyadic-interval-converging-sequence*:
assumes $t \in \{0..T\} \wedge T \neq 0$
shows $\exists s. \forall n. s \ n \in \text{dyadic-interval } 0 \ T - \{t\} \wedge s \longrightarrow t$
 $\langle \text{proof} \rangle$

lemma *dyadic-interval-dense*: $\text{closure } (\text{dyadic-interval } 0 \ T) = \{0..T\}$
 $\langle \text{proof} \rangle$

corollary *dyadic-interval-islimpt*:
assumes $T > 0 \wedge t \in \{0..T\}$
shows $t \text{ islimpt dyadic-interval } 0 \ T$
 $\langle \text{proof} \rangle$

corollary *at-within-dyadic-interval-nontrivial*[simp]:
assumes $T > 0 \wedge t \in \{0..T\}$
shows $(\text{at } t \text{ within dyadic-interval } 0 \ T) \neq \text{bot}$
 $\langle \text{proof} \rangle$

lemma *dyadic-interval-step-finite*[simp]: *finite* (*dyadic-interval-step* *n* *S* *T*)
 ⟨*proof*⟩

lemma *dyadic-interval-countable*[simp]: *countable* (*dyadic-interval* *S* *T*)
 ⟨*proof*⟩

lemma *floor-pow2-add-leq*:
 fixes *T* :: *real*
 shows $\lfloor 2^n * T \rfloor / 2^n \leq \lfloor 2^{n+k} * T \rfloor / 2^{n+k}$
 ⟨*proof*⟩

corollary *floor-pow2-mono*: *mono* ($\lambda n. \lfloor 2^n * (T :: \text{real}) \rfloor / 2^n$)
 ⟨*proof*⟩

lemma *dyadic-interval-step-Max*: $T \geq 0 \implies \text{Max} (\text{dyadic-interval-step } n \ 0 \ T) = \lfloor 2^n * T \rfloor / 2^n$
 ⟨*proof*⟩

lemma *dyadic-interval-step-subset*:
 $n \leq m \implies \text{dyadic-interval-step } n \ 0 \ T \subseteq \text{dyadic-interval-step } m \ 0 \ T$
 ⟨*proof*⟩

corollary *dyadic-interval-step-mono*:
 assumes $x \in \text{dyadic-interval-step } n \ 0 \ T$ $n \leq m$
 shows $x \in \text{dyadic-interval-step } m \ 0 \ T$
 ⟨*proof*⟩

lemma *dyadic-as-natural*:
 assumes $x \in \text{dyadic-interval-step } n \ 0 \ T$
 shows $\exists! k. x = \text{real } k / 2^n$
 ⟨*proof*⟩

lemma *dyadic-of-natural*:
 assumes $\text{real } k / 2^n \leq T$
 shows $\text{real } k / 2^n \in \text{dyadic-interval-step } n \ 0 \ T$
 ⟨*proof*⟩

lemma *dyadic-interval-minus*:
 assumes $x \in \text{dyadic-interval-step } n \ 0 \ T$ $y \in \text{dyadic-interval-step } n \ 0 \ T$ $x \leq y$
 shows $y - x \in \text{dyadic-interval-step } n \ 0 \ T$
 ⟨*proof*⟩

lemma *dyadic-times-nat*: $x \in \text{dyadic-interval-step } n \ 0 \ T \implies (x * 2^n) \in \mathbb{N}$
 ⟨*proof*⟩

definition *dyadic-expansion* $x \ n \ b \ k \equiv \text{set } b \subseteq \{0,1\}$
 $\wedge \text{length } b = n \wedge x = \text{real-of-int } k + (\sum_{m \in \{1..n\}} \text{real } (b ! (m-1)) / 2^m)$

lemma *dyadic-expansionI*:

assumes $set\ b \subseteq \{0,1\}$ $length\ b = n$ $x = k + (\sum_{m \in \{1..n\}} (b\ !\ (m-1)) / 2^m)$
shows $dyadic-expansion\ x\ n\ b\ k$
 $\langle proof \rangle$

lemma *dyadic-expansionD*:
assumes $dyadic-expansion\ x\ n\ b\ k$
shows $set\ b \subseteq \{0,1\}$
and $length\ b = n$
and $x = k + (\sum_{m \in \{1..n\}} (b\ !\ (m-1)) / 2^m)$
 $\langle proof \rangle$

lemma *dyadic-expansion-ex*:
assumes $x \in dyadic-interval-step\ n\ 0\ T$
shows $\exists b\ k. dyadic-expansion\ x\ n\ b\ k$
 $\langle proof \rangle$

lemma *dyadic-expansion-frac-le-1*:
assumes $dyadic-expansion\ x\ n\ b\ k$
shows $(\sum_{m \in \{1..n\}} (b\ !\ (m-1)) / 2^m) < 1$
 $\langle proof \rangle$

lemma *dyadic-expansion-frac-range*:
assumes $dyadic-expansion\ x\ n\ b\ k\ m \in \{1..n\}$
shows $b\ !\ (m-1) \in \{0,1\}$
 $\langle proof \rangle$

lemma *dyadic-expansion-interval*:
assumes $dyadic-expansion\ x\ n\ b\ k\ x \in \{S..T\}$
shows $x \in dyadic-interval-step\ n\ S\ T$
 $\langle proof \rangle$

lemma *dyadic-expansion-nth-geq*:
assumes $dyadic-expansion\ x\ n\ b\ k\ m \in \{1..n\}\ b\ !\ (m-1) = 1$
shows $x \geq k + 1/2^m$
 $\langle proof \rangle$

lemma *dyadic-expansion-frac-geq-0*:
assumes $dyadic-expansion\ x\ n\ b\ k$
shows $(\sum_{m \in \{1..n\}} (b\ !\ (m-1)) / 2^m) \geq 0$
 $\langle proof \rangle$

lemma *dyadic-expansion-frac*:
assumes $dyadic-expansion\ x\ n\ b\ k$
shows $frac\ x = (\sum_{m \in \{1..n\}} (b\ !\ (m-1)) / 2^m)$
 $\langle proof \rangle$

lemma *dyadic-expansion-floor*:
assumes $dyadic-expansion\ x\ n\ b\ k$

shows $k = \lfloor x \rfloor$
 $\langle \text{proof} \rangle$

lemma *sum-interval-pow2-inv*: $(\sum m \in \{ \text{Suc } l..n \}. (1 :: \text{real}) / 2^{\wedge} m) = 1 / 2^{\wedge} l - 1 / 2^{\wedge} n$ **if** $l < n$
 $\langle \text{proof} \rangle$

lemma *dyadic-expansion-unique*:
assumes *dyadic-expansion* x n b k
and *dyadic-expansion* x n c j
shows $b = c \wedge j = k$
 $\langle \text{proof} \rangle$

end

3 Hölder continuity

theory *Holder-Continuous*
imports *HOL-Analysis.Analysis*
begin

Hölder continuity is a weaker version of Lipschitz continuity.

definition *holder-at-within* :: $\text{real} \Rightarrow 'a \text{ set} \Rightarrow 'a \Rightarrow ('a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space}) \Rightarrow \text{bool}$ **where**
 $\text{holder-at-within } \gamma \ D \ r \ \varphi \equiv \gamma \in \{0 <..1\} \wedge$
 $(\exists \varepsilon > 0. \exists C \geq 0. \forall s \in D. \text{dist } r \ s < \varepsilon \longrightarrow \text{dist } (\varphi \ r) (\varphi \ s) \leq C * \text{dist } r \ s^{\text{powr } \gamma})$

definition *local-holder-on* :: $\text{real} \Rightarrow 'a :: \text{metric-space} \text{ set} \Rightarrow ('a \Rightarrow 'b :: \text{metric-space}) \Rightarrow \text{bool}$ **where**
 $\text{local-holder-on } \gamma \ D \ \varphi \equiv \gamma \in \{0 <..1\} \wedge$
 $(\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in D. \forall s \in D. \text{dist } s \ t < \varepsilon \wedge \text{dist } r \ t < \varepsilon \longrightarrow \text{dist } (\varphi \ r) (\varphi \ s) \leq C * \text{dist } r \ s^{\text{powr } \gamma}))$

definition *holder-on* :: $\text{real} \Rightarrow 'a :: \text{metric-space} \text{ set} \Rightarrow ('a \Rightarrow 'b :: \text{metric-space}) \Rightarrow \text{bool}$ $(--\text{holder'-on } 1000)$ **where**
 $\gamma\text{-holder-on } D \ \varphi \longleftrightarrow \gamma \in \{0 <..1\} \wedge (\exists C \geq 0. (\forall r \in D. \forall s \in D. \text{dist } (\varphi \ r) (\varphi \ s) \leq C * \text{dist } r \ s^{\text{powr } \gamma}))$

lemma *holder-onI*:
assumes $\gamma \in \{0 <..1\} \exists C \geq 0. (\forall r \in D. \forall s \in D. \text{dist } (\varphi \ r) (\varphi \ s) \leq C * \text{dist } r \ s^{\text{powr } \gamma})$
shows $\gamma\text{-holder-on } D \ \varphi$
 $\langle \text{proof} \rangle$

We prove various equivalent formulations of local holder continuity, using open and closed balls and inequalities.

lemma *local-holder-on-cball*:

$local\text{-}holder\text{-}on\ \gamma\ D\ \varphi \longleftrightarrow \gamma \in \{0 < .. 1\} \wedge$
 $(\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in cball\ t\ \varepsilon \cap D. \forall s \in cball\ t\ \varepsilon \cap D. dist\ (\varphi\ r)\ (\varphi\ s)$
 $\leq C * dist\ r\ s\ powr\ \gamma))$
 $(is\ ?L \longleftrightarrow ?R)$
 $\langle proof \rangle$

corollary *local-holder-on-leq-def*: $local\text{-}holder\text{-}on\ \gamma\ D\ \varphi \longleftrightarrow \gamma \in \{0 < .. 1\} \wedge$
 $(\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in D. \forall s \in D. dist\ s\ t \leq \varepsilon \wedge dist\ r\ t \leq \varepsilon \longrightarrow dist\ (\varphi\ r)$
 $)\ (\varphi\ s) \leq C * dist\ r\ s\ powr\ \gamma))$
 $\langle proof \rangle$

corollary *local-holder-on-ball*: $local\text{-}holder\text{-}on\ \gamma\ D\ \varphi \longleftrightarrow \gamma \in \{0 < .. 1\} \wedge$
 $(\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in ball\ t\ \varepsilon \cap D. \forall s \in ball\ t\ \varepsilon \cap D. dist\ (\varphi\ r)\ (\varphi\ s)$
 $\leq C * dist\ r\ s\ powr\ \gamma))$
 $\langle proof \rangle$

lemma *local-holder-on-altdef*:
assumes $D \neq \{\}$
shows $local\text{-}holder\text{-}on\ \gamma\ D\ \varphi = (\forall t \in D. (\exists \varepsilon > 0. (\gamma\text{-}holder\text{-}on\ ((cball\ t\ \varepsilon) \cap D)\ \varphi)))$
 $\langle proof \rangle$

lemma *local-holder-on-cong*[*cong*]:
assumes $\gamma = \varepsilon\ C = D\ \bigwedge x. x \in C \implies \varphi\ x = \psi\ x$
shows $local\text{-}holder\text{-}on\ \gamma\ C\ \varphi \longleftrightarrow local\text{-}holder\text{-}on\ \varepsilon\ D\ \psi$
 $\langle proof \rangle$

lemma *local-holder-onI*:
assumes $\gamma \in \{0 < .. 1\}\ (\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in D. \forall s \in D. dist\ s\ t < \varepsilon \wedge$
 $dist\ r\ t < \varepsilon \longrightarrow dist\ (\varphi\ r)\ (\varphi\ s) \leq C * dist\ r\ s\ powr\ \gamma))$
shows $local\text{-}holder\text{-}on\ \gamma\ D\ \varphi$
 $\langle proof \rangle$

lemma *local-holder-ballI*:
assumes $\gamma \in \{0 < .. 1\}$
and $\bigwedge t. t \in D \implies \exists \varepsilon > 0. \exists C \geq 0. \forall r \in ball\ t\ \varepsilon \cap D. \forall s \in ball\ t\ \varepsilon \cap D.$
 $dist\ (\varphi\ r)\ (\varphi\ s) \leq C * dist\ r\ s\ powr\ \gamma$
shows $local\text{-}holder\text{-}on\ \gamma\ D\ \varphi$
 $\langle proof \rangle$

lemma *local-holder-onE*:
assumes *local-holder*: $local\text{-}holder\text{-}on\ \gamma\ D\ \varphi$
and *gamma*: $\gamma \in \{0 < .. 1\}$
and $t \in D$
obtains $\varepsilon\ C$ **where** $\varepsilon > 0\ C \geq 0$
 $\bigwedge r\ s. r \in ball\ t\ \varepsilon \cap D \implies s \in ball\ t\ \varepsilon \cap D \implies dist\ (\varphi\ r)\ (\varphi\ s) \leq C * dist\ r$
 $s\ powr\ \gamma$
 $\langle proof \rangle$

Holder continuity matches up with the existing definitions in *HOL-Analysis.Lipschitz*

lemma *holder-1-eq-lipschitz*: $1\text{-holder-on } D \varphi = (\exists C. \text{lipschitz-on } C D \varphi)$
 ⟨proof⟩

lemma *local-holder-1-eq-local-lipschitz*:
 assumes $T \neq \{\}$
 shows $\text{local-holder-on } 1 D \varphi = \text{local-lipschitz } T D (\lambda\cdot. \varphi)$
 ⟨proof⟩

lemma *local-holder-refine*:
 assumes $g: \text{local-holder-on } g D \varphi$ and $h: h \leq g$ and $h > 0$
 shows $\text{local-holder-on } h D \varphi$
 ⟨proof⟩

lemma *holder-uniform-continuous*:
 assumes $\gamma\text{-holder-on } X \varphi$
 shows $\text{uniformly-continuous-on } X \varphi$
 ⟨proof⟩

corollary *holder-on-continuous-on*: $\gamma\text{-holder-on } X \varphi \implies \text{continuous-on } X \varphi$
 ⟨proof⟩

lemma *holder-implies-local-holder*: $\gamma\text{-holder-on } D \varphi \implies \text{local-holder-on } \gamma D \varphi$
 ⟨proof⟩

lemma *local-holder-imp-continuous*:
 assumes $\text{local-holder: local-holder-on } \gamma X \varphi$
 shows $\text{continuous-on } X \varphi$
 ⟨proof⟩

lemma *local-holder-compact-imp-holder*:
 assumes $\text{compact } I$ and $\text{local-holder-on } \gamma I \varphi$
 shows $\gamma\text{-holder-on } I \varphi$
 ⟨proof⟩

lemma *holder-const*: $\gamma\text{-holder-on } C (\lambda\cdot. c) \longleftrightarrow \gamma \in \{0 < .. 1\}$
 ⟨proof⟩

lemma *local-holder-const*: $\text{local-holder-on } \gamma C (\lambda\cdot. c) \longleftrightarrow \gamma \in \{0 < .. 1\}$
 ⟨proof⟩

end

4 Convergence in measure

theory *Measure-Convergence*
 imports *HOL-Probability.Probability*
 begin

We use `measure` rather than `emeasure` because `ennreal` is not a metric space, which we need to reason about convergence. By intersecting with the set of finite measure `A`, we don't run into issues where infinity is collapsed to 0. For finite measures this definition is equal to the definition without set `A` – see below.

definition `tendsto-measure` :: `'b measure  $\Rightarrow$  ('a  $\Rightarrow$  'b  $\Rightarrow$  ('c :: {second-countable-topology, metric-space}))  $\Rightarrow$  ('b  $\Rightarrow$  'c)  $\Rightarrow$  'a filter  $\Rightarrow$  bool`
where `tendsto-measure M X l F $\equiv$ ($\forall n. X\ n \in \text{borel-measurable } M$) $\wedge$ $l \in \text{borel-measurable } M \wedge$
 $(\forall \varepsilon > 0. \forall A \in \text{fmeasurable } M.$
 $((\lambda n. \text{measure } M (\{\omega \in \text{space } M. \text{dist } (X\ n\ \omega) (l\ \omega) > \varepsilon\} \cap A)) \longrightarrow 0) F)$`

abbreviation (in `prob-space`) `tendsto-prob` (**infixr** `$\longrightarrow_P$` 55) **where**
`(f  $\longrightarrow_P$  l) F  $\equiv$  tendsto-measure M f l F`

lemma `tendsto-measure-measurable[measurable-dest]`:
`tendsto-measure M X l F  $\implies$  X n  $\in$  borel-measurable M`
`<proof>`

lemma `tendsto-measure-measurable-lim[measurable-dest]`:
`tendsto-measure M X l F  $\implies$  l  $\in$  borel-measurable M`
`<proof>`

lemma `tendsto-measure-mono`: `F  $\leq$  F'  $\implies$  tendsto-measure M f l F'  $\implies$  tendsto-measure M f l F`
`<proof>`

lemma `tendsto-measureI`:
assumes `[measurable]:  $\bigwedge n. X\ n \in \text{borel-measurable } M$   $l \in \text{borel-measurable } M$`
and `$\bigwedge \varepsilon A. \varepsilon > 0 \implies A \in \text{fmeasurable } M \implies$`
 `$((\lambda n. \text{measure } M (\{\omega \in \text{space } M. \text{dist } (X\ n\ \omega) (l\ \omega) > \varepsilon\} \cap A)) \longrightarrow 0) F$`
shows `tendsto-measure M X l F`
`<proof>`

lemma (in `finite-measure`) `finite-tendsto-measureI`:
assumes `[measurable]:  $\bigwedge n. f'\ n \in \text{borel-measurable } M$   $f \in \text{borel-measurable } M$`
and `$\bigwedge \varepsilon. \varepsilon > 0 \implies ((\lambda n. \text{measure } M (\{\omega \in \text{space } M. \text{dist } (f'\ n\ \omega) (f\ \omega) > \varepsilon\}) \longrightarrow 0) F$`
shows `tendsto-measure M f' f F`
`<proof>`

lemma (in `finite-measure`) `finite-tendsto-measureD`:
assumes `[measurable]: tendsto-measure M f' f F`
shows `$(\forall \varepsilon > 0. ((\lambda n. \text{measure } M (\{\omega \in \text{space } M. \text{dist } (f'\ n\ \omega) (f\ \omega) > \varepsilon\}) \longrightarrow 0) F)$`
`<proof>`

lemma (in `finite-measure`) `tendsto-measure-leq`:

assumes $[measurable]: \bigwedge n. f' n \in \text{borel-measurable } M \ f \in \text{borel-measurable } M$
shows $\text{tendsto-measure } M \ f' \ f \ F \longleftrightarrow$
 $(\forall \varepsilon > 0. ((\lambda n. \text{measure } M \ \{\omega \in \text{space } M. \text{dist } (f' n \ \omega) \ (f \ \omega) \geq \varepsilon\}) \longrightarrow 0)$
 $F) \text{ (is } ?L \longleftrightarrow ?R)$
 $\langle \text{proof} \rangle$

abbreviation $\text{LIMSEQ-measure } M \ f \ l \equiv \text{tendsto-measure } M \ f \ l \text{ sequentially}$

lemma $\text{LIMSEQ-measure-def}: \text{LIMSEQ-measure } M \ f \ l \longleftrightarrow$
 $(\forall n. f \ n \in \text{borel-measurable } M) \wedge (l \in \text{borel-measurable } M) \wedge$
 $(\forall \varepsilon > 0. \forall A \in \text{fmeasurable } M.$
 $(\lambda n. \text{measure } M \ (\{\omega \in \text{space } M. \text{dist } (f \ n \ \omega) \ (l \ \omega) > \varepsilon\} \cap A)) \longrightarrow 0)$
 $\langle \text{proof} \rangle$

lemma LIMSEQ-measureD :
assumes $\text{LIMSEQ-measure } M \ f \ l \ \varepsilon > 0 \ A \in \text{fmeasurable } M$
shows $(\lambda n. \text{measure } M \ (\{\omega \in \text{space } M. \text{dist } (f \ n \ \omega) \ (l \ \omega) > \varepsilon\} \cap A)) \longrightarrow 0$
 $\langle \text{proof} \rangle$

lemma fmeasurable-inter : $\llbracket A \in \text{sets } M; B \in \text{fmeasurable } M \rrbracket \implies A \cap B \in \text{fmeasurable } M$
 $\langle \text{proof} \rangle$

lemma $\text{LIMSEQ-measure-emeasure}$:
assumes $\text{LIMSEQ-measure } M \ f \ l \ \varepsilon > 0 \ A \in \text{fmeasurable } M$
and $[measurable]: \bigwedge i. f \ i \in \text{borel-measurable } M \ l \in \text{borel-measurable } M$
shows $(\lambda n. \text{emeasure } M \ (\{\omega \in \text{space } M. \text{dist } (f \ n \ \omega) \ (l \ \omega) > \varepsilon\} \cap A)) \longrightarrow 0$
 $\langle \text{proof} \rangle$

lemma $\text{measure-Lim-within-LIMSEQ}$:
fixes $a :: 'a :: \text{first-countable-topology}$
assumes $\bigwedge t. X \ t \in \text{borel-measurable } M \ L \in \text{borel-measurable } M$
assumes $\bigwedge S. \llbracket (\forall n. S \ n \neq a \wedge S \ n \in T); S \longrightarrow a \rrbracket \implies \text{LIMSEQ-measure } M$
 $(\lambda n. X \ (S \ n)) \ L$
shows $\text{tendsto-measure } M \ X \ L \ (\text{at } a \text{ within } T)$
 $\langle \text{proof} \rangle$

definition $\text{tendsto-AE} :: 'b \text{ measure} \Rightarrow ('a \Rightarrow 'b \Rightarrow 'c :: \text{topological-space}) \Rightarrow ('b$
 $\Rightarrow 'c) \Rightarrow 'a \text{ filter} \Rightarrow \text{bool}$ **where**
 $\text{tendsto-AE } M \ f' \ l \ F \longleftrightarrow (AE \ \omega \text{ in } M. ((\lambda n. f' n \ \omega) \longrightarrow l \ \omega) \ F)$

lemma $\text{LIMSEQ-ae-pointwise}$: $(\bigwedge x. (\lambda n. f \ n \ x) \longrightarrow l \ x) \implies \text{tendsto-AE } M \ f \ l$
 sequentially
 $\langle \text{proof} \rangle$

lemma $\text{tendsto-AE-within-LIMSEQ}$:
fixes $a :: 'a :: \text{first-countable-topology}$
assumes $\bigwedge S. \llbracket (\forall n. S \ n \neq a \wedge S \ n \in T); S \longrightarrow a \rrbracket \implies \text{tendsto-AE } M \ (\lambda n.$
 $X \ (S \ n)) \ L \text{ sequentially}$

shows *tendsto-AE* M X L (at a within T)
 ⟨proof⟩

lemma *LIMSEQ-dominated-convergence*:
fixes $X :: \text{nat} \Rightarrow \text{real}$
assumes $X \longrightarrow L$ ($\bigwedge n. Y\ n \leq X\ n$) ($\bigwedge n. Y\ n \geq L$)
shows $Y \longrightarrow L$
 ⟨proof⟩

Klenke remark 6.4

lemma *measure-conv-imp-AE-sequentially*:
assumes [*measurable*]: $\bigwedge n. f'\ n \in \text{borel-measurable } M$ $f \in \text{borel-measurable } M$
and *tendsto-AE* M f' f *sequentially*
shows *LIMSEQ-measure* M f' f
 ⟨proof⟩

corollary *LIMSEQ-measure-pointwise*:
assumes $\bigwedge x. (\bigwedge n. f\ n\ x) \longrightarrow f'\ x$ $\bigwedge n. f\ n \in \text{borel-measurable } M$ $f' \in \text{borel-measurable } M$
shows *LIMSEQ-measure* M f f'
 ⟨proof⟩

lemma *Lim-measure-pointwise*:
fixes $a :: 'a :: \text{first-countable-topology}$
assumes $\bigwedge x. ((\bigwedge n. f\ n\ x) \longrightarrow f'\ x) (\text{at } a \text{ within } T)$ $\bigwedge n. f\ n \in \text{borel-measurable } M$ $f' \in \text{borel-measurable } M$
shows *tendsto-measure* M f f' (at a within T)
 ⟨proof⟩

corollary *measure-conv-imp-AE-at-within*:
fixes $x :: 'a :: \text{first-countable-topology}$
assumes [*measurable*]: $\bigwedge n. f'\ n \in \text{borel-measurable } M$ $f \in \text{borel-measurable } M$
and *tendsto-AE* M f' f (at x within S)
shows *tendsto-measure* M f' f (at x within S)
 ⟨proof⟩

Klenke remark 6.5

lemma (in *sigma-finite-measure*) *LIMSEQ-measure-unique-AE*:
fixes $f :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{metric-space}\}$
assumes [*measurable*]: $\bigwedge n. f\ n \in \text{borel-measurable } M$ $l \in \text{borel-measurable } M$ $l' \in \text{borel-measurable } M$
and *LIMSEQ-measure* M f l *LIMSEQ-measure* M f l'
shows *AE* x in $M. l\ x = l'\ x$
 ⟨proof⟩

corollary (in *sigma-finite-measure*) *LIMSEQ-ae-unique-AE*:
fixes $f :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{metric-space}\}$
assumes $\bigwedge n. f\ n \in \text{borel-measurable } M$ $l \in \text{borel-measurable } M$ $l' \in \text{borel-measurable } M$
shows *AE* x in $M. l\ x = l'\ x$

and *tendsto-AE* $M f l$ *sequentially tendsto-AE* $M f l'$ *sequentially*
shows *AE* x *in* M . $l x = l' x$
 <proof>

lemma (*in sigma-finite-measure*) *tendsto-measure-at-within-eq-AE*:
fixes $f :: 'b :: \text{first-countable-topology} \Rightarrow 'a \Rightarrow 'c :: \{\text{second-countable-topology}, \text{metric-space}\}$
assumes *f-measurable*: $\bigwedge x. x \in S \implies f x \in \text{borel-measurable } M$
and *l-measurable*: $l \in \text{borel-measurable } M \ l' \in \text{borel-measurable } M$
and *tendsto*: *tendsto-measure* $M f l$ (at t *within* S) *tendsto-measure* $M f l'$ (at t *within* S)
and *not-bot*: (at t *within* S) $\neq \perp$
shows *AE* x *in* M . $l x = l' x$
 <proof>
end

5 Stochastic processes

theory *Stochastic-Processes*
imports *Kolmogorov-Chentsov-Extras Dyadic-Interval*
begin

A stochastic process is an indexed collection of random variables. For compatibility with **product_prob_space** we don't enforce conditions on the index set I in the assumptions.

locale *stochastic-process* = *prob-space* +
fixes $M' :: 'b \text{ measure}$
and $I :: 't \text{ set}$
and $X :: 't \Rightarrow 'a \Rightarrow 'b$
assumes *random-process[measurable]*: $\bigwedge i. \text{random-variable } M' (X i)$

sublocale *stochastic-process* $\subseteq$ *product*: *product-prob-space* ($\lambda t. \text{distr } M M' (X t)$)
 <proof>

lemma (*in prob-space*) *stochastic-process* I :
assumes $\bigwedge i. \text{random-variable } M' (X i)$
shows *stochastic-process* $M M' X$
 <proof>

typedef ($'t, 'a, 'b$) *stochastic-process* =
 $\{(M :: 'a \text{ measure}, M' :: 'b \text{ measure}, I :: 't \text{ set}, X :: 't \Rightarrow 'a \Rightarrow 'b).$
 $\text{stochastic-process } M M' X\}$
 <proof>

setup-lifting *type-definition-stochastic-process*

lift-definition *proc-source* :: ($'t, 'a, 'b$) *stochastic-process* $\Rightarrow 'a \text{ measure}$
is *fst* <proof>

interpretation *proc-source*: *prob-space* *proc-source* *X*
 ⟨*proof*⟩

lift-definition *proc-target* :: ('t,'a,'b) *stochastic-process* ⇒ 'b *measure*
 is *fst* ∘ *snd* ⟨*proof*⟩

lift-definition *proc-index* :: ('t,'a,'b) *stochastic-process* ⇒ 't *set*
 is *fst* ∘ *snd* ∘ *snd* ⟨*proof*⟩

lift-definition *process* :: ('t,'a,'b) *stochastic-process* ⇒ 't ⇒ 'a ⇒ 'b
 is *snd* ∘ *snd* ∘ *snd* ⟨*proof*⟩

declare [[*coercion process*]]

lemma *stochastic-process-construct* [*simp*]: *stochastic-process* (*proc-source* *X*) (*proc-target* *X*) (*process* *X*)
 ⟨*proof*⟩

interpretation *stochastic-process* *proc-source* *X* *proc-target* *X* *proc-index* *X* *process* *X*
 ⟨*proof*⟩

lemma *stochastic-process-measurable* [*measurable*]: *process* *X* *t* ∈ (*proc-source* *X*)
 →_{*M*} (*proc-target* *X*)
 ⟨*proof*⟩

Here we construct a process on a given index set. For this we need to produce measurable functions for indices outside the index set; we use the constant function, but it needs to point at an element of the target set to be measurable.

context *prob-space*
begin

lift-definition *process-of* :: 'b *measure* ⇒ 't *set* ⇒ ('t ⇒ 'a ⇒ 'b) ⇒ 'b ⇒ ('t,'a,'b) *stochastic-process*
 is λ *M'* *I* *X* ω. if (∀ *t* ∈ *I*. *X* *t* ∈ *M* →_{*M*} *M'*) ∧ ω ∈ *space* *M'*
 then (*M*, *M'*, *I*, (λ*t*. if *t* ∈ *I* then *X* *t* else (λ-. ω)))
 else (return (*sigma* *UNIV* { {}, *UNIV* }) (*SOME* *x*. *True*), *sigma* *UNIV* *UNIV*,
I, λ- -. ω)
 ⟨*proof*⟩

lemma *index-process-of* [*simp*]: *proc-index* (*process-of* *M'* *I* *X* ω) = *I*
 ⟨*proof*⟩

lemma
assumes ∀ *t* ∈ *I*. *X* *t* ∈ *M* →_{*M*} *M'* ω ∈ *space* *M'*
shows
source-process-of [*simp*]: *proc-source* (*process-of* *M'* *I* *X* ω) = *M* **and**
target-process-of [*simp*]: *proc-target* (*process-of* *M'* *I* *X* ω) = *M'* **and**

process-process-of[simp]: process (process-of M' I X ω) = (λt . if $t \in I$ then $X t$ else (λ -. ω))
<proof>

lemma *process-of-apply:*

assumes $\forall t \in I. X t \in M \rightarrow_M M' \omega \in \text{space } M' t \in I$

shows *process (process-of M' I X ω) t = X t*

<proof>

end

We define the finite-dimensional distributions of our process.

lift-definition *distributions :: ('t, 'a, 'b) stochastic-process $\Rightarrow$ 't set $\Rightarrow$ ('t $\Rightarrow$ 'b) measure*

is $\lambda(M, M', -, X) T. (\Pi_M t \in T. \text{distr } M M' (X t))$ *<proof>*

lemma *distributions-altdef: distributions X T = ($\Pi_M t \in T. \text{distr (proc-source X) (proc-target X) (X t)$)*

<proof>

lemma *prob-space-distributions: prob-space (distributions X J)*

<proof>

lemma *sets-distributions: sets (distributions X J) = sets (PiM J (λ -. (proc-target X)))*

<proof>

lemma *space-distributions: space (distributions X J) = ($\Pi_E i \in J. \text{space (proc-target X)}$)*

<proof>

lemma *emeasure-distributions:*

assumes *finite J $\bigwedge j. j \in J \implies A j \in \text{sets (proc-target X)}$*

shows *emeasure (distributions X J) (PiE J A) = ($\prod j \in J. \text{emeasure (distr (proc-source X) (proc-target X) (X j)) (A j)}$)*

<proof>

interpretation *projective-family (proc-index X) distributions X (λ -. proc-target X)*

<proof>

locale *polish-stochastic = stochastic-process M borel :: 'b::polish-space measure I X for M and I and X*

lemma *distributed-cong-random-variable:*

assumes $M = K N = L A E x \text{ in } M. X x = Y x X \in M \rightarrow_M N Y \in K \rightarrow_M L$
 $f \in \text{borel-measurable } N$

shows *distributed M N X f $\longleftrightarrow$ distributed K L Y f*

$\langle \text{proof} \rangle$

For all sorted lists of indices, the increments specified by this list are independent

lift-definition *indep-increments* :: ($'t :: \text{linorder}, 'a, 'b :: \text{minus}$) *stochastic-process* $\Rightarrow \text{bool}$ **is**

$\lambda(M, M', I, X).$
 $(\forall l. \text{set } l \subseteq I \wedge \text{sorted } l \wedge \text{length } l \geq 2 \longrightarrow$
 $\text{prob-space.indep-vars } M (\lambda-. M') (\lambda k v. X (!k) v - X (! (k-1)) v) \{1..<\text{length } l\})$ $\langle \text{proof} \rangle$

lemma *indep-incrementsE*:

assumes *indep-increments* X
and $\text{set } l \subseteq \text{proc-index } X \wedge \text{sorted } l \wedge \text{length } l \geq 2$
shows $\text{prob-space.indep-vars } (\text{proc-source } X) (\lambda-. \text{proc-target } X)$
 $(\lambda k v. X (!k) v - X (! (k-1)) v) \{1..<\text{length } l\}$
 $\langle \text{proof} \rangle$

lemma *indep-incrementsI*:

assumes $\bigwedge l. \text{set } l \subseteq \text{proc-index } X \Longrightarrow \text{sorted } l \Longrightarrow \text{length } l \geq 2 \Longrightarrow$
 $\text{prob-space.indep-vars } (\text{proc-source } X) (\lambda-. \text{proc-target } X) (\lambda k v. X (!k) v - X$
 $(!(k-1)) v) \{1..<\text{length } l\}$
shows *indep-increments* X
 $\langle \text{proof} \rangle$

lemma *indep-increments-indep-var*:

assumes *indep-increments* X $h \in \text{proc-index } X$ $j \in \text{proc-index } X$ $k \in \text{proc-index } X$ $h \leq j$ $j \leq k$
shows $\text{prob-space.indep-var } (\text{proc-source } X) (\text{proc-target } X) (\lambda v. X j v - X h v)$
 $(\text{proc-target } X) (\lambda v. X k v - X j v)$
 $\langle \text{proof} \rangle$

definition *stationary-increments* $X \longleftrightarrow (\forall t1 t2 k. t1 > 0 \wedge t2 > 0 \wedge k > 0 \longrightarrow$

$\text{distr } (\text{proc-source } X) (\text{proc-target } X) (\lambda v. X (t1 + k) v - X t1 v) =$
 $\text{distr } (\text{proc-source } X) (\text{proc-target } X) (\lambda v. X (t2 + k) v - X t2 v))$

Processes on the same source measure space, with the same index space, but not necessarily the same target measure since we only care about the measurable target space, not the measure

lift-definition *compatible* :: ($'t, 'a, 'b$) *stochastic-process* $\Rightarrow ('t, 'a, 'b)$ *stochastic-process* $\Rightarrow \text{bool}$

is $\lambda(Mx, M'x, Ix, X) (My, M'y, Iy, -). Mx = My \wedge \text{sets } M'x = \text{sets } M'y \wedge Ix = Iy$ $\langle \text{proof} \rangle$

lemma *compatibleI*:

assumes $\text{proc-source } X = \text{proc-source } Y$ $\text{sets } (\text{proc-target } X) = \text{sets } (\text{proc-target } Y)$
 $\text{proc-index } X = \text{proc-index } Y$

shows *compatible* $X Y$
 $\langle \text{proof} \rangle$

lemma

assumes *compatible* $X Y$

shows

compatible-source $[dest]$: *proc-source* $X = \text{proc-source } Y$ **and**
compatible-target $[dest]$: *sets* (*proc-target* X) = *sets* (*proc-target* Y) **and**
compatible-index $[dest]$: *proc-index* $X = \text{proc-index } Y$

$\langle \text{proof} \rangle$

lemma *compatible-refl* $[simp]$: *compatible* $X X$

$\langle \text{proof} \rangle$

lemma *compatible-sym*: *compatible* $X Y \implies \text{compatible } Y X$

$\langle \text{proof} \rangle$

lemma *compatible-trans*:

assumes *compatible* $X Y$ *compatible* $Y Z$

shows *compatible* $X Z$

$\langle \text{proof} \rangle$

lemma (**in** *prob-space*) *compatible-process-of*:

assumes *measurable*: $\forall t \in I. X t \in M \rightarrow_M M' \forall t \in I. Y t \in M \rightarrow_M M'$

and $a \in \text{space } M' b \in \text{space } M'$

shows *compatible* (*process-of* $M' I X a$) (*process-of* $M' I Y b$)

$\langle \text{proof} \rangle$

definition *modification* :: $('t, 'a, 'b)$ *stochastic-process* $\Rightarrow ('t, 'a, 'b)$ *stochastic-process*

$\Rightarrow \text{bool}$ **where**

modification $X Y \longleftrightarrow \text{compatible } X Y \wedge (\forall t \in \text{proc-index } X. \text{AE } x \text{ in } \text{proc-source } X. X t x = Y t x)$

lemma *modificationI* $[intro]$:

assumes *compatible* $X Y \bigwedge t. t \in \text{proc-index } X \implies \text{AE } x \text{ in } \text{proc-source } X. X t x = Y t x$

shows *modification* $X Y$

$\langle \text{proof} \rangle$

lemma *modificationD* $[dest]$:

assumes *modification* $X Y$

shows *compatible* $X Y$

and $\bigwedge t. t \in \text{proc-index } X \implies \text{AE } x \text{ in } \text{proc-source } X. X t x = Y t x$

$\langle \text{proof} \rangle$

lemma *modification-null-set*:

assumes *modification* $X Y t \in \text{proc-index } X$

obtains N **where** $\{x \in \text{space } (\text{proc-source } X). X t x \neq Y t x\} \subseteq N N \in \text{null-sets } (\text{proc-source } X)$

$\langle \text{proof} \rangle$

lemma *modification-refl* [simp]: *modification* X X
 $\langle \text{proof} \rangle$

lemma *modification-sym*: *modification* X $Y \implies$ *modification* Y X
 $\langle \text{proof} \rangle$

lemma *modification-trans*:
 assumes *modification* X Y *modification* Y Z
 shows *modification* X Z
 $\langle \text{proof} \rangle$

lemma *modification-imp-identical-distributions*:
 assumes *modification*: *modification* X Y
 and *index*: $T \subseteq \text{proc-index } X$
 shows *distributions* X $T =$ *distributions* Y T
 $\langle \text{proof} \rangle$

definition *indistinguishable* :: $(t, 'a, 'b)$ *stochastic-process* $\Rightarrow$ $(t, 'a, 'b)$ *stochastic-process*
 $\Rightarrow$ *bool* **where**
 indistinguishable X $Y \iff$ *compatible* X $Y \wedge$
 $(\exists N \in \text{null-sets } (\text{proc-source } X). \forall t \in \text{proc-index } X. \{x \in \text{space } (\text{proc-source } X). \\ X\ t\ x \neq Y\ t\ x\} \subseteq N)$

lemma *indistinguishableI*:
 assumes *compatible* X Y
 and $\exists N \in \text{null-sets } (\text{proc-source } X). (\forall t \in \text{proc-index } X. \{x \in \text{space } (\text{proc-source } X). \\ X\ t\ x \neq Y\ t\ x\} \subseteq N)$
 shows *indistinguishable* X Y
 $\langle \text{proof} \rangle$

lemma *indistinguishable-null-set*:
 assumes *indistinguishable* X Y
 obtains N **where**
 $N \in \text{null-sets } (\text{proc-source } X)$
 $\bigwedge t. t \in \text{proc-index } X \implies \{x \in \text{space } (\text{proc-source } X). X\ t\ x \neq Y\ t\ x\} \subseteq N$
 $\langle \text{proof} \rangle$

lemma *indistinguishableD*:
 assumes *indistinguishable* X Y
 shows *compatible* X Y
 and $\exists N \in \text{null-sets } (\text{proc-source } X). (\forall t \in \text{proc-index } X. \{x \in \text{space } (\text{proc-source } X). \\ X\ t\ x \neq Y\ t\ x\} \subseteq N)$
 $\langle \text{proof} \rangle$

lemma *indistinguishable-eq-AE*:
 assumes *indistinguishable* X Y
 shows *AE* x *in* *proc-source* $X. \forall t \in \text{proc-index } X. X\ t\ x = Y\ t\ x$

<proof>

lemma *indistinguishable-null-ex:*

assumes *indistinguishable X Y*

shows $\exists N \in \text{null-sets } (\text{proc-source } X). \{x \in \text{space } (\text{proc-source } X). \exists t \in \text{proc-index } X. X \ t \ x \neq Y \ t \ x\} \subseteq N$

<proof>

lemma *indistinguishable-refl [simp]: indistinguishable X X*

<proof>

lemma *indistinguishable-sym: indistinguishable X Y $\implies$ indistinguishable Y X*

<proof>

lemma *indistinguishable-trans:*

assumes *indistinguishable X Y indistinguishable Y Z*

shows *indistinguishable X Z*

<proof>

lemma *indistinguishable-modification: indistinguishable X Y $\implies$ modification X Y*

<proof>

Klenke 21.5(i)

lemma *modification-countable:*

assumes *modification X Y countable (proc-index X)*

shows *indistinguishable X Y*

<proof>

Klenke 21.5(ii). The textbook statement is more general - we reduce right continuity to regular continuity

lemma *modification-continuous-indistinguishable:*

fixes *X :: (real, 'a, 'b :: metric-space) stochastic-process*

assumes *modification: modification X Y*

and interval: $\exists T > 0. \text{proc-index } X = \{0..T\}$

and rc: *AE ω in proc-source X. continuous-on (proc-index X) ($\lambda t. X \ t \ \omega$)*

(is AE ω in proc-source X. ?cont-X ω)

AE ω in proc-source Y. continuous-on (proc-index Y) ($\lambda t. Y \ t \ \omega$)

(is AE ω in proc-source Y. ?cont-Y ω)

shows *indistinguishable X Y*

<proof>

lift-definition *restrict-index :: ('t, 'a, 'b) stochastic-process $\Rightarrow$ 't set $\Rightarrow$ ('t, 'a, 'b) stochastic-process*

is $\lambda(M, M', I, X) \ T. (M, M', T, X) \ \langle \text{proof} \rangle$

lemma

shows

$\text{restrict-index-source}[\text{simp}]: \text{proc-source } (\text{restrict-index } X \ T) = \text{proc-source } X$
and
 $\text{restrict-index-target}[\text{simp}]: \text{proc-target } (\text{restrict-index } X \ T) = \text{proc-target } X$ **and**
 $\text{restrict-index-index}[\text{simp}]: \text{proc-index } (\text{restrict-index } X \ T) = T$ **and**
 $\text{restrict-index-process}[\text{simp}]: \text{process } (\text{restrict-index } X \ T) = \text{process } X$
 $\langle \text{proof} \rangle$

lemma $\text{restrict-index-override}[\text{simp}]: \text{restrict-index } (\text{restrict-index } X \ T) \ S = \text{restrict-index } X \ S$
 $\langle \text{proof} \rangle$

lemma $\text{compatible-restrict-index}$:
assumes $\text{compatible } X \ Y$
shows $\text{compatible } (\text{restrict-index } X \ S) \ (\text{restrict-index } Y \ S)$
 $\langle \text{proof} \rangle$

lemma $\text{modification-restrict-index}$:
assumes $\text{modification } X \ Y \ S \subseteq \text{proc-index } X$
shows $\text{modification } (\text{restrict-index } X \ S) \ (\text{restrict-index } Y \ S)$
 $\langle \text{proof} \rangle$

lemma $\text{indistinguishable-restrict-index}$:
assumes $\text{indistinguishable } X \ Y \ S \subseteq \text{proc-index } X$
shows $\text{indistinguishable } (\text{restrict-index } X \ S) \ (\text{restrict-index } Y \ S)$
 $\langle \text{proof} \rangle$

lemma $\text{AE-eq-minus} [\text{intro}]$:
fixes $a :: 'a \Rightarrow ('b :: \text{real-normed-vector})$
assumes $\text{AE } x \text{ in } M. a \ x = b \ x \ \text{AE } x \text{ in } M. c \ x = d \ x$
shows $\text{AE } x \text{ in } M. a \ x - c \ x = b \ x - d \ x$
 $\langle \text{proof} \rangle$

lemma $\text{modification-indep-increments}$:
fixes $X \ Y :: ('a :: \text{linorder}, 'b, 'c :: \{\text{second-countable-topology}, \text{real-normed-vector}\})$
 $\text{stochastic-process}$
assumes $\text{modification } X \ Y \ \text{sets } (\text{proc-target } Y) = \text{sets borel}$
shows $\text{indep-increments } X \Longrightarrow \text{indep-increments } Y$
 $\langle \text{proof} \rangle$

end

6 The Kolomgorov-Chentsov theorem

theory $\text{Kolmogorov-Chentsov}$
imports $\text{Stochastic-Processes Holder-Continuous Dyadic-Interval Measure-Convergence}$
begin

6.1 Supporting lemmas

The main contribution of this file is the Kolmogorov-Chentsov theorem: given a stochastic process that satisfies some continuity properties, we can construct a Hölder continuous modification. We first prove some auxiliary lemmas before moving on to the main construction.

Klenke 5.11: Markov inequality. Compare with $\llbracket (\lambda x. ?u\ x * \text{indicator } ?A\ x) \in \text{borel-measurable } ?M; ?A \in \text{sets } ?M \rrbracket \implies \text{emeasure } ?M \{x \in ?A. 1 \leq ?c * ?u\ x\} \leq ?c * \text{set-nn-integral } ?M\ ?A\ ?u$

lemma *nn-integral-Markov-inequality-extended:*

fixes $f :: \text{real} \Rightarrow \text{ennreal}$ **and** $\varepsilon :: \text{real}$ **and** $X :: 'a \Rightarrow \text{real}$
assumes *mono: mono-on (range $X \cup \{0<..\}$) f*
and *finite: $\bigwedge x. f\ x < \infty$*
and $e: \varepsilon > 0\ f\ \varepsilon > 0$
and *[measurable]: $X \in \text{borel-measurable } M$*
shows $\text{emeasure } M \{p \in \text{space } M. (X\ p) \geq \varepsilon\} \leq (\int^+ x. f\ (X\ x)\ \partial M) / f\ \varepsilon$
<proof>

lemma *nn-integral-Markov-inequality-extended-rnv:*

fixes $f :: \text{real} \Rightarrow \text{real}$ **and** $\varepsilon :: \text{real}$ **and** $X :: 'a \Rightarrow 'b :: \text{real-normed-vector}$
assumes *[measurable]: $X \in \text{borel-measurable } M$*
and *mono: mono-on $\{0<..\}$ f*
and $e: \varepsilon > 0\ f\ \varepsilon > 0$
shows $\text{emeasure } M \{p \in \text{space } M. \text{norm } (X\ p) \geq \varepsilon\} \leq (\int^+ x. f\ (\text{norm } (X\ x))\ \partial M) / f\ \varepsilon$
<proof>

6.2 Kolmogorov-Chentsov

Klenke theorem 21.6 - Kolmogorov-Chentsov

locale *Kolmogorov-Chentsov =*

fixes $X :: (\text{real}, 'a, 'b :: \text{polish-space})\ \text{stochastic-process}$
and $a\ b\ C\ \gamma :: \text{real}$
assumes *index[simp]: proc-index $X = \{0<..\}$*
and *target-borel[simp]: proc-target $X = \text{borel}$*
and *gt-0: $a > 0\ b > 0\ C > 0$*
and *b-leq-a: $b \leq a$*
and *gamma: $\gamma \in \{0<.. $b/a\}$$*
and *expectation: $\bigwedge s\ t. \llbracket s \geq 0; t \geq 0 \rrbracket \implies$*
 $(\int^+ x. \text{dist } (X\ t\ x)\ (X\ s\ x)\ \text{powr } a\ \partial \text{proc-source } X) \leq C * \text{dist } t\ s\ \text{powr } (1+b)$
begin

lemma *gamma-0-1[simp]: $\gamma \in \{0<.. $1\}$$*
<proof>

lemma *gamma-gt-0[simp]: $\gamma > 0$*

$\langle \text{proof} \rangle$

lemma *gamma-le-1[simp]*: $\gamma \leq 1$
 $\langle \text{proof} \rangle$

abbreviation *source* $\equiv$ *proc-source* *X*

lemma *X-borel-measurable[measurable]*: $X \ t \in \text{borel-measurable source}$ **for** *t*
 $\langle \text{proof} \rangle$

lemma *markov*: $\mathcal{P}(x \text{ in source. } \varepsilon \leq \text{dist } (X \ t \ x) \ (X \ s \ x)) \leq (C * \text{dist } t \ s \ \text{powr } (1 + b)) / \varepsilon \ \text{powr } a$
if $s \geq 0 \ t \geq 0 \ \varepsilon > 0$ **for** $s \ t \ \varepsilon$
 $\langle \text{proof} \rangle$

lemma *conv-in-prob*:
assumes $t \geq 0$
shows *tendsto-measure (proc-source X) X (X t) (at t within {0..})*
 $\langle \text{proof} \rangle$

lemma *conv-in-prob-finite*:
assumes $t \geq 0$
shows *tendsto-measure (proc-source X) X (X t) (at t within {0..T})*
 $\langle \text{proof} \rangle$

lemma *incr*: $\mathcal{P}(x \text{ in source. } 2 \ \text{powr } (-\gamma * n) \leq \text{dist } (X \ ((k - 1) * 2 \ \text{powr } - n) \ x) \ (X \ (k * 2 \ \text{powr } - n) \ x))$
 $\leq C * 2 \ \text{powr } (-n * (1 + b - a * \gamma))$
if $k \geq 1 \ n \geq 0$ **for** $k \ n$
 $\langle \text{proof} \rangle$

end

In order to construct the modification of X , it suffices to construct a modification of X on $\{0..T\}$ for all finite T , from which we construct the modification on $\{0..\}$ via a countable union.

locale *Kolmogorov-Chentsov-finite* = *Kolmogorov-Chentsov* +
fixes $T :: \text{real}$
assumes *zero-le-T*: $0 < T$
begin

A_n will characterise the set of states with increments that exceed the bounds required for Hölder continuity. As $n \rightarrow \infty$, this approaches the set of states for which X is not Hölder continuous. We define N as this limit, and show that N is a null set. On $\omega \in \Omega - N$, we show that $X(\omega)$ is Hölder continuous (and therefore uniformly continuous) on the dyadic rationals, and construct a modification by taking the continuous extension on the reals.

definition $A \equiv \lambda n. \text{if } 2^{-n} * T < 1 \text{ then space source else}$

$\{x \in \text{space source}.$
 $\text{Max } \{ \text{dist } (X \text{ (real-of-int } (k - 1) * 2^{\text{powr} - \text{real } n}) x) (X \text{ (real-of-int } k * 2^{\text{powr} - \text{real } n}) x)$
 $\mid k. k \in \{1..[2^n * T]\} \} \geq 2^{\text{powr} - \gamma * \text{real } n} \}$

abbreviation $B \equiv \lambda n. (\bigcup m. A (m + n))$

abbreviation $N \equiv \bigcap (\text{range } B)$

lemma $A\text{-geq}$: $2^n * T \geq 1 \implies A n = \{x \in \text{space source}.$
 $\text{Max } \{ \text{dist } (X \text{ (real-of-int } (k - 1) * 2^{\text{powr} - \text{real } n}) x) (X \text{ (real-of-int } k * 2^{\text{powr} - \text{real } n}) x)$
 $\mid k. k \in \{1..[2^n * T]\} \} \geq 2^{\text{powr} - \gamma * \text{real } n} \}$ **for** n
 $\langle \text{proof} \rangle$

lemma $A\text{-measurable}[\text{measurable}]$: $A n \in \text{sets source}$
 $\langle \text{proof} \rangle$

lemma $\text{emeasure-}A\text{-leq}$:
fixes $n :: \text{nat}$
assumes $[\text{simp}]$: $2^n * T \geq 1$
shows $\text{emeasure source } (A n) \leq C * T * 2^{\text{powr} - n * (b - a * \gamma)}$
 $\langle \text{proof} \rangle$

lemma $\text{measure-}A\text{-leq}$:
assumes $2^n * T \geq 1$
shows $\text{measure source } (A n) \leq C * T * 2^{\text{powr} - n * (b - a * \gamma)}$
 $\langle \text{proof} \rangle$

lemma $\text{summable-}A$: $\text{summable } (\lambda m. \text{measure source } (A m))$
 $\langle \text{proof} \rangle$

lemma $\text{lim-}B$: $(\lambda n. \text{measure source } (B n)) \longrightarrow 0$
 $\langle \text{proof} \rangle$

lemma $N\text{-null}$: $N \in \text{null-sets source}$
 $\langle \text{proof} \rangle$

lemma $\text{notin-}N\text{-index}$:
assumes $\omega \in \text{space source} - N$
obtains n_0 **where** $\omega \notin (\bigcup n. A (n + n_0))$
 $\langle \text{proof} \rangle$

context
fixes ω
assumes $\omega: \omega \in \text{space source} - N$
begin

definition $n_0 \equiv \text{SOME } m. \omega \notin (\bigcup n. A (n + m)) \wedge m > 0$

lemma

shows $n\text{-zero}$: $\omega \notin (\bigcup n. A (n + n_0))$
and $n\text{-zero-nonzero}$: $n_0 > 0$
 $\langle \text{proof} \rangle$

lemma $n\text{zero-ge}$: $\bigwedge n. n \geq n_0 \implies 2^{\wedge n} * T \geq 1$
 $\langle \text{proof} \rangle$

lemma ωnotin : $\bigwedge n. n \geq n_0 \implies \omega \notin A n$
 $\langle \text{proof} \rangle$

Klenke 21.7

lemma $X\text{-dyadic-incr}$:

assumes $k \in \{1.. \lfloor 2^{\wedge n} * T \rfloor\}$ $n \geq n_0$
shows $\text{dist } (X ((\text{real-of-int } k-1)/2^{\wedge n}) \omega) (X (\text{real-of-int } k/2^{\wedge n}) \omega) < 2^{\text{powr}} (-\gamma * n)$
 $\langle \text{proof} \rangle$

Klenke (21.8)

lemma dist-dyadic-mn :

assumes mn : $n_0 \leq n$ $n \leq m$
and $t\text{-dyadic}$: $t \in \text{dyadic-interval-step } m \ 0 \ T$
and $u\text{-dyadic-n}$: $u \in \text{dyadic-interval-step } n \ 0 \ T$
and ut : $u \leq t$ $t - u < 2/2^{\wedge n}$
shows $\text{dist } (X u \omega) (X t \omega) \leq 2^{\text{powr}} (-\gamma * n) / (1 - 2^{\text{powr}} - \gamma)$
 $\langle \text{proof} \rangle$

lemma dist-dyadic-fixed :

assumes mn : $n_0 \leq n$ $n \leq m$
and $s\text{-dyadic}$: $s \in \text{dyadic-interval-step } m \ 0 \ T$
and $t\text{-dyadic}$: $t \in \text{dyadic-interval-step } m \ 0 \ T$
and st : $s \leq t$ $t - s \leq 1/2^{\wedge n}$
shows $\text{dist } (X t \omega) (X s \omega) \leq 2 * 2^{\text{powr}} (-\gamma * n) / (1 - 2^{\text{powr}} - \gamma)$
 $\langle \text{proof} \rangle$

definition $C_0 \equiv 2 * 2^{\text{powr}} \gamma / (1 - 2^{\text{powr}} - \gamma)$

lemma $C\text{-zero-ge}[\text{simp}]$: $C_0 > 0$
 $\langle \text{proof} \rangle$

Klenke (21.9)

Let $s, t \in D$ with $|s - t| \leq \frac{1}{2_0^n}$. By choosing the minimal $n \geq n_0$ such that $|t - s| \geq 2^{-n}$, we obtain by $\llbracket n_0 \leq ?n; ?n \leq ?m; ?s \in \text{dyadic-interval-step } ?m \ 0 \ T; ?t \in \text{dyadic-interval-step } ?m \ 0 \ T; ?s \leq ?t; ?t - ?s \leq 1 / 2^{?n} \rrbracket \implies \text{dist } (\text{process } X \ ?t \ \omega) (\text{process } X \ ?s \ \omega) \leq 2 * 2^{\text{powr}} (-\gamma * \text{real } ?n) / (1 - 2^{\text{powr}} - \gamma)$:

$$|X_t(\omega) - X_s(\omega)| \leq C_0 |t - s|^\gamma$$

lemma *dist-dyadic*:

assumes $t: t \in \text{dyadic-interval } 0 \ T$

and $s: s \in \text{dyadic-interval } 0 \ T$

and $st\text{-}dist: \text{dist } t \ s \leq 1 / 2^{\wedge n_0}$

shows $\text{dist } (X \ t \ \omega) \ (X \ s \ \omega) \leq C_0 * (\text{dist } t \ s)^{\text{powr } \gamma}$

<proof>

definition $K \equiv C_0 * (2^{\wedge nat} \lceil 2^{\wedge n_0} * T \rceil)^{\text{powr } (1 - \gamma)}$

lemma $C_0\text{-le-}K$: $C_0 \leq K$

<proof>

lemma $K\text{-pos}$: $0 < K$

<proof>

Klenke (21.10)

lemma $X\text{-dyadic-le-}K'$:

assumes $dyadic: s \in \text{dyadic-interval } 0 \ T \ t \in \text{dyadic-interval } 0 \ T$

and $st: s \leq t$

shows $\text{dist } (X \ s \ \omega) \ (X \ t \ \omega) \leq K * \text{dist } s \ t^{\text{powr } \gamma}$

<proof>

lemma $X\text{-dyadic-le-}K$:

assumes $s \in \text{dyadic-interval } 0 \ T$

and $t \in \text{dyadic-interval } 0 \ T$

shows $\text{dist } (X \ s \ \omega) \ (X \ t \ \omega) \leq K * \text{dist } s \ t^{\text{powr } \gamma}$

<proof>

corollary $holder\text{-}dyadic$: $\gamma\text{-holder-on } (\text{dyadic-interval } 0 \ T) \ (\lambda t. X \ t \ \omega)$

<proof>

lemma $uniformly\text{-continuous-dyadic}$: $uniformly\text{-continuous-on } (\text{dyadic-interval } 0 \ T) \ (\lambda t. X \ t \ \omega)$

<proof>

lemma $Lim\text{-exists}$: $\exists L. ((\lambda s. X \ s \ \omega) \longrightarrow L) \text{ (at } t \text{ within } (\text{dyadic-interval } 0 \ T))$

if $t \in \{0..T\}$

<proof>

lemma $Lim\text{-unique}$: $\exists! L. ((\lambda s. X \ s \ \omega) \longrightarrow L) \text{ (at } t \text{ within } (\text{dyadic-interval } 0 \ T))$

if $t \in \{0..T\}$

<proof>

definition $L \equiv (\lambda t. (Lim \text{ (at } t \text{ within } \text{dyadic-interval } 0 \ T) \ (\lambda s. X \ s \ \omega)))$

lemma $X\text{-tendsto-}L$:

assumes $t \in \{0..T\}$

shows $((\lambda s. X s \omega) \longrightarrow L t) \text{ (at } t \text{ within (dyadic-interval } 0 T))$
 $\langle \text{proof} \rangle$

lemma *L-dist-K*:
assumes $s: s \in \{0..T\}$
and $t: t \in \{0..T\}$
shows $\text{dist } (L s) (L t) \leq K * \text{dist } s t \text{ powr } \gamma$
 $\langle \text{proof} \rangle$

corollary *L-holder: γ -holder-on $\{0..T\}$ L*
 $\langle \text{proof} \rangle$

corollary *L-local-holder: local-holder-on $\gamma \{0..T\}$ L*
 $\langle \text{proof} \rangle$

lemma *X-dyadic-eq-L*:
assumes $t \in \text{dyadic-interval } 0 T$
shows $X t \omega = L t$
 $\langle \text{proof} \rangle$
end

definition *default* :: 'b **where** *default* = (SOME x. True)

definition *X-tilde* :: real $\Rightarrow$ 'a $\Rightarrow$ 'b **where**
 $X\text{-tilde} \equiv (\lambda t \omega. \text{if } \omega \in N \text{ then default else } (Lim \text{ (at } t \text{ within dyadic-interval } 0 T) (\lambda s. X s \omega)))$

lemma *X-tilde-not-N-Lim*:
assumes $\omega \in \text{space source} - N$
shows $X\text{-tilde } t \omega = Lim \text{ (at } t \text{ within dyadic-interval } 0 T) (\lambda s. X s \omega)$
 $\langle \text{proof} \rangle$

lemma *X-tilde-not-N-L*:
assumes $\omega \in \text{space source} - N$
shows $X\text{-tilde } t \omega = L \omega t$
 $\langle \text{proof} \rangle$

lemma *local-holder-X-tilde: local-holder-on $\gamma \{0..T\}$ $(\lambda t. X\text{-tilde } t \omega)$*
if $\omega \in \text{space source}$ **for** ω
 $\langle \text{proof} \rangle$

corollary *X-tilde-eq-L-AE: AE ω in source. $X\text{-tilde } t \omega = L \omega t$*
 $\langle \text{proof} \rangle$

corollary *X-tilde-eq-Lim-AE*:
 $AE \omega \text{ in source. } X\text{-tilde } t \omega = Lim \text{ (at } t \text{ within dyadic-interval } 0 T) (\lambda s. X s \omega)$
 $\langle \text{proof} \rangle$

lemma *X-tilde-tendsto-AE: $t \in \{0..T\} \implies \text{tendsto-AE source } X (X\text{-tilde } t) \text{ (at } t$*

within dyadic-interval 0 T)
 ⟨proof⟩

end

context Kolmogorov-Chentsov-finite
begin

By (21.5) $0 \leq ?t \implies \text{tendsto-measure source (process } X) \text{ (process } X ?t) \text{ (at } ?t \text{ within } \{0..?T\})}$ and (21.11) $? \omega \in \text{space source} - (\bigcap_n \bigcup_m A (m + n)) \implies L ? \omega \equiv \lambda t. \text{Lim (at } t \text{ within dyadic-interval } 0 T) (\lambda s. \text{process } X s ? \omega), P[X \neq \tilde{X}] = 0$

lemma *X-tilde-measurable[measurable]:*
assumes $t \in \{0..T\}$
shows $X\text{-tilde } t \in \text{borel-measurable source}$
 ⟨proof⟩

lemma *X-eq-X-tilde-AE: AE ω in source. $X t \omega = X\text{-tilde } t \omega$ if $t \in \{0..T\}$ for t*
 ⟨proof⟩

lemma *X-tilde-modification: modification (restrict-index X $\{0..T\}$) (prob-space.process-of source (proc-target X) $\{0..T\}$ X-tilde default)*
 ⟨proof⟩

end

We have now shown that we can construct a modification of X for any interval $\{0..T\}$. We want to extend this result to construct a modification on the interval $\{0..\}$ - this can be constructed by gluing together all modifications with natural-valued T which results in a countable union of modifications, which itself is a modification.

context Kolmogorov-Chentsov
begin

lemma *Kolmogorov-Chentsov-finite: $T > 0 \implies \text{Kolmogorov-Chentsov-finite } X a b C \gamma T$*
 ⟨proof⟩

definition $\text{Mod} \equiv \lambda T. \text{SOME } Y. \text{modification (restrict-index } X \{0..T\}) Y \wedge (\forall x \in \text{space source. local-holder-on } \gamma \{0..T\} (\lambda t. Y t x))$

lemma *Mod: modification (restrict-index X $\{0..T\}$) (Mod T)*
 $(\forall x \in \text{space source. local-holder-on } \gamma \{0..T\} (\lambda t. (\text{Mod } T) t x))$ **if** $0 < T$ **for** T
 ⟨proof⟩

lemma *compatible-Mod: compatible (restrict-index X $\{0..T\}$) (Mod T) if $0 < T$ for T*
 ⟨proof⟩

lemma *Mod-source[simp]*: $\text{proc-source } (\text{Mod } T) = \text{source}$ **if** $0 < T$ **for** T
 $\langle \text{proof} \rangle$

lemma *Mod-target*: $\text{sets } (\text{proc-target } (\text{Mod } T)) = \text{sets } (\text{proc-target } X)$ **if** $0 < T$ **for** T
 $\langle \text{proof} \rangle$

lemma *Mod-index[simp]*: $0 < T \implies \text{proc-index } (\text{Mod } T) = \{0..T\}$
 $\langle \text{proof} \rangle$

lemma *indistinguishable-mod*:
 $\text{indistinguishable } (\text{restrict-index } (\text{Mod } S) \{0..\min S T\}) (\text{restrict-index } (\text{Mod } T) \{0..\min S T\})$
if $S > 0$ $T > 0$ **for** $S T$
 $\langle \text{proof} \rangle$

definition $N S T \equiv \text{SOME } N. N \in \text{null-sets source} \wedge \{\omega \in \text{space source}. \exists t \in \{0..\min S T\}. (\text{Mod } S) t \omega \neq (\text{Mod } T) t \omega\} \subseteq N$

lemma N :
assumes $S > 0$ $T > 0$
shows $N S T \in \text{null-sets source} \wedge \{\omega \in \text{space source}. \exists t \in \{0..\min S T\}. (\text{Mod } S) t \omega \neq (\text{Mod } T) t \omega\} \subseteq N S T$
 $\langle \text{proof} \rangle$

definition $N\text{-inf}$ **where** $N\text{-inf} \equiv (\bigcup S \in \mathbb{N} - \{0\}. (\bigcup T \in \mathbb{N} - \{0\}. N S T))$

lemma $N\text{-inf-null}$: $N\text{-inf} \in \text{null-sets source}$
 $\langle \text{proof} \rangle$

lemma *Mod-eq-N-inf*: $(\text{Mod } S) t \omega = (\text{Mod } T) t \omega$
if $\omega \in \text{space source} - N\text{-inf } t \in \{0..\min S T\}$ $S \in \mathbb{N} - \{0\}$ $T \in \mathbb{N} - \{0\}$ **for** $\omega t S T$
 $\langle \text{proof} \rangle$

definition $\text{default} :: 'b$ **where** $\text{default} = (\text{SOME } x. \text{True})$

definition $X\text{-mod} \equiv \lambda t \omega. \text{if } \omega \in \text{space source} - N\text{-inf} \text{ then } (\text{Mod } \lfloor t+1 \rfloor) t \omega \text{ else } \text{default}$

definition $X\text{-mod-process} \equiv \text{prob-space.process-of source } (\text{proc-target } X) \{0..\} X\text{-mod default}$

lemma *Mod-measurable[measurable]*: $\forall t \in \{0..\}. X\text{-mod } t \in \text{source} \rightarrow_M \text{proc-target } X$
 $\langle \text{proof} \rangle$

lemma *modification-X-mod-process*: $\text{modification } X X\text{-mod-process}$
 $\langle \text{proof} \rangle$

lemma *local-holder-X-mod*: *local-holder-on* $\gamma \{0..\}$ $(\lambda t. X\text{-mod } t \ \omega)$ **for** ω
 $\langle \text{proof} \rangle$

lemma *local-holder-X-mod-process*: *local-holder-on* $\gamma \{0..\}$ $(\lambda t. X\text{-mod-process } t \ \omega)$
for ω
 $\langle \text{proof} \rangle$

theorem *continuous-modification*:
 $\exists X'. \text{modification } X \ X' \wedge (\forall \omega. \text{local-holder-on } \gamma \{0..\} (\lambda t. X' \ t \ \omega))$
 $\langle \text{proof} \rangle$

end

theorem *Kolmogorov-Chentsov*:
fixes $X :: (\text{real}, 'a, 'b :: \text{polish-space}) \text{ stochastic-process}$
and $a \ b \ C \ \gamma :: \text{real}$
assumes *index[simp]*: *proc-index* $X = \{0..\}$
and *target-borel[simp]*: *proc-target* $X = \text{borel}$
and *gt-0*: $a > 0 \ b > 0 \ C > 0$
and *b-leq-a*: $b \leq a$
and *gamma*: $\gamma \in \{0 < .. < b/a\}$
and *expectation*: $\bigwedge s \ t. \llbracket s \geq 0; t \geq 0 \rrbracket \implies$
 $(\int^+ x. \text{dist } (X \ t \ x) \ (X \ s \ x) \ \text{powr } a \ \partial \text{proc-source } X) \leq C * \text{dist } t \ s \ \text{powr } (1+b)$
shows $\exists X'. \text{modification } X \ X' \wedge (\forall \omega. \text{local-holder-on } \gamma \{0..\} (\lambda t. X' \ t \ \omega))$
 $\langle \text{proof} \rangle$
end

References

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