

Khovanskii's Theorem

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Abstract

We formalise the proof of an important theorem in additive combinatorics due to Khovanskii [2, 3], attesting that the cardinality of the set of all sums of n many elements of A , where A is a finite subset of an abelian group, is a polynomial in n for all sufficiently large n . We follow a proof of the theorem due to Nathanson and Ruzsa [4, 5] as presented in the notes “Introduction to Additive Combinatorics” by Timothy Gowers [1] for the University of Cambridge.

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1 Product Operator for Commutative Monoids

theory *FiniteProduct*

imports

Jacobson-Basic-Algebra.Group-Theory

begin

1.1 Products over Finite Sets

context *commutative-monoid* **begin**

definition *M-ify* $x \equiv \text{if } x \in M \text{ then } x \text{ else } \mathbf{1}$

definition *fincomp* $f A \equiv \text{if finite } A \text{ then } \text{Finite-Set.fold } (\lambda x y. f x \cdot M\text{-ify } y) \mathbf{1} A \text{ else } \mathbf{1}$

lemma *fincomp-empty* [*simp*]: $\text{fincomp } f \ \{\} = \mathbf{1}$
<proof>

lemma *fincomp-infinite*[*simp*]: $\text{infinite } A \implies \text{fincomp } f A = \mathbf{1}$
<proof>

lemma *left-commute*: $\llbracket a \in M; b \in M; c \in M \rrbracket \implies b \cdot (a \cdot c) = a \cdot (b \cdot c)$
<proof>

lemma *comp-fun-commute-onI*:

assumes $f \in F \rightarrow M$

shows *comp-fun-commute-on* $F (\lambda x y. f x \cdot M\text{-ify } y)$

<proof>

lemma *fincomp-closed* [*simp*]:

assumes $f \in F \rightarrow M$

shows $\text{fincomp } f F \in M$

<proof>

lemma *fincomp-insert* [*simp*]:

assumes $F: \text{finite } F \ a \notin F$ **and** $f: f \in F \rightarrow M \ f a \in M$

shows $\text{fincomp } f (\text{insert } a F) = f a \cdot \text{fincomp } f F$

<proof>

lemma *fincomp-unit-eqI*: $(\bigwedge x. x \in A \implies f x = \mathbf{1}) \implies \text{fincomp } f A = \mathbf{1}$

<proof>

lemma *fincomp-unit* [*simp*]: $\text{fincomp } (\lambda i. \mathbf{1}) A = \mathbf{1}$

<proof>

lemma *funcset-Int-left* [*simp, intro*]:

$\llbracket f \in A \rightarrow C; g \in B \rightarrow C \rrbracket \implies f \in A \ \text{Int } B \rightarrow C$

$\langle proof \rangle$

lemma *funcset-Un-left* [iff]:

$$(f \in A \text{ Un } B \rightarrow C) = (f \in A \rightarrow C \wedge f \in B \rightarrow C)$$

$\langle proof \rangle$

lemma *fincomp-Un-Int*:

$$\llbracket \text{finite } A; \text{ finite } B; g \in A \rightarrow M; g \in B \rightarrow M \rrbracket \implies$$

$$\text{fincomp } g (A \cup B) \cdot \text{fincomp } g (A \cap B) =$$

$$\text{fincomp } g A \cdot \text{fincomp } g B$$

— The reversed orientation looks more natural, but LOOPS as a simprule!

$\langle proof \rangle$

lemma *fincomp-Un-disjoint*:

$$\llbracket \text{finite } A; \text{ finite } B; A \cap B = \{\}; g \in A \rightarrow M; g \in B \rightarrow M \rrbracket$$

$$\implies \text{fincomp } g (A \cup B) = \text{fincomp } g A \cdot \text{fincomp } g B$$

$\langle proof \rangle$

lemma *fincomp-comp*:

$$\llbracket f \in A \rightarrow M; g \in A \rightarrow M \rrbracket \implies \text{fincomp } (\lambda x. f x \cdot g x) A = (\text{fincomp } f A \cdot \text{fincomp } g A)$$

$\langle proof \rangle$

lemma *fincomp-cong'*:

$$\text{assumes } A = B \text{ } g \in B \rightarrow M \wedge i. i \in B \implies f i = g i$$

$$\text{shows } \text{fincomp } f A = \text{fincomp } g B$$

$\langle proof \rangle$

lemma *fincomp-cong*:

$$\text{assumes } A = B \text{ } g \in B \rightarrow M \wedge i. i \in B = \text{simp} \implies f i = g i$$

$$\text{shows } \text{fincomp } f A = \text{fincomp } g B$$

$\langle proof \rangle$

Usually, if this rule causes a failed congruence proof error, the reason is that the premise $g \in B \rightarrow M$ cannot be shown. Adding *Pi-def* to the simpset is often useful. For this reason, *fincomp-cong* is not added to the simpset by default.

lemma *fincomp-0* [simp]:

$$f \in \{0 :: \text{nat}\} \rightarrow M \implies \text{fincomp } f \{..0\} = f 0$$

$\langle proof \rangle$

lemma *fincomp-0'*: $f \in \{..n\} \rightarrow M \implies (f 0) \cdot \text{fincomp } f \{\text{Suc } 0..n\} = \text{fincomp } f \{..n\}$

$\langle proof \rangle$

lemma *fincomp-Suc* [simp]:

$$f \in \{..\text{Suc } n\} \rightarrow M \implies \text{fincomp } f \{..\text{Suc } n\} = (f (\text{Suc } n)) \cdot \text{fincomp } f \{..n\}$$

$\langle proof \rangle$

lemma *fincomp-Suc2*:

$f \in \{..Suc\ n\} \rightarrow M \implies fincomp\ f\ \{..Suc\ n\} = (fincomp\ (\%i.\ f\ (Suc\ i))\ \{..n\} \cdot f\ 0)$
<proof>

lemma *fincomp-Suc3*:

assumes $f \in \{..n\ :: nat\} \rightarrow M$
shows $fincomp\ f\ \{..n\} = (f\ n) \cdot fincomp\ f\ \{..<\ n\}$
<proof>

lemma *fincomp-reindex*:

$f \in (h\ 'A) \rightarrow M \implies$
 $inj\text{-on}\ h\ A \implies fincomp\ f\ (h\ 'A) = fincomp\ (\lambda x.\ f\ (h\ x))\ A$
<proof>

lemma *fincomp-const*:

assumes $a\ [simp]: a \in M$
shows $fincomp\ (\lambda x.\ a)\ A = rec\text{-nat}\ \mathbf{1}\ (\lambda u.\ (\cdot)\ a)\ (card\ A)$
<proof>

lemma *fincomp-singleton*:

assumes $i\text{-in}\ A: i \in A$ **and** $fin\text{-}A: finite\ A$ **and** $f\text{-}Pi: f \in A \rightarrow M$
shows $fincomp\ (\lambda j.\ if\ i = j\ then\ f\ j\ else\ \mathbf{1})\ A = f\ i$
<proof>

lemma *fincomp-singleton-swap*:

assumes $i\text{-in}\ A: i \in A$ **and** $fin\text{-}A: finite\ A$ **and** $f\text{-}Pi: f \in A \rightarrow M$
shows $fincomp\ (\lambda j.\ if\ j = i\ then\ f\ j\ else\ \mathbf{1})\ A = f\ i$
<proof>

lemma *fincomp-mono-neutral-cong-left*:

assumes $finite\ B$
and $A \subseteq B$
and $1: \bigwedge i.\ i \in B - A \implies h\ i = \mathbf{1}$
and $gh: \bigwedge x.\ x \in A \implies g\ x = h\ x$
and $h: h \in B \rightarrow M$
shows $fincomp\ g\ A = fincomp\ h\ B$
<proof>

lemma *fincomp-mono-neutral-cong-right*:

assumes $finite\ B$
and $A \subseteq B \bigwedge i.\ i \in B - A \implies g\ i = \mathbf{1} \bigwedge x.\ x \in A \implies g\ x = h\ x$ $g \in B \rightarrow M$
shows $fincomp\ g\ B = fincomp\ h\ A$
<proof>

lemma *fincomp-mono-neutral-cong*:

assumes $[simp]: finite\ B\ finite\ A$
and $*$: $\bigwedge i.\ i \in B - A \implies h\ i = \mathbf{1} \bigwedge i.\ i \in A - B \implies g\ i = \mathbf{1}$
and $gh: \bigwedge x.\ x \in A \cap B \implies g\ x = h\ x$

and $g: g \in A \rightarrow M$
and $h: h \in B \rightarrow M$
shows $\text{fincomp } g \ A = \text{fincomp } h \ B$
 $\langle \text{proof} \rangle$

lemma *fincomp-UN-disjoint*:

assumes

$\text{finite } I \wedge i. i \in I \implies \text{finite } (A \ i) \text{ pairwise } (\lambda i \ j. \text{disjnt } (A \ i) \ (A \ j)) \ I$

$\wedge i \ x. i \in I \implies x \in A \ i \implies g \ x \in M$

shows $\text{fincomp } g \ (\bigcup (A \ ' I)) = \text{fincomp } (\lambda i. \text{fincomp } g \ (A \ i)) \ I$

$\langle \text{proof} \rangle$

lemma *fincomp-Union-disjoint*:

$\llbracket \text{finite } C; \wedge A. A \in C \implies \text{finite } A \wedge (\forall x \in A. f \ x \in M); \text{pairwise disjoint } C \rrbracket \implies$

$\text{fincomp } f \ (\bigcup C) = \text{fincomp } (\text{fincomp } f) \ C$

$\langle \text{proof} \rangle$

end

1.2 Results for Abelian Groups

context *abelian-group* **begin**

lemma *fincomp-inverse*:

$f \in A \rightarrow G \implies \text{fincomp } (\lambda x. \text{inverse } (f \ x)) \ A = \text{inverse } (\text{fincomp } f \ A)$

$\langle \text{proof} \rangle$

Jeremy Avigad. This should be generalized to arbitrary groups, not just Abelian ones, using Lagrange's theorem.

lemma *power-order-eq-one*:

assumes *fin [simp]*: $\text{finite } G$

and *a [simp]*: $a \in G$

shows $\text{rec-nat } \mathbf{1} \ (\lambda u. (\cdot) \ a) \ (\text{card } G) = \mathbf{1}$

$\langle \text{proof} \rangle$

end

end

2 Khovanskii's Theorem

We formalise the proof of an important theorem in additive combinatorics due to Khovanskii, attesting that the cardinality of the set of all sums of n many elements of A , where A is a finite subset of an abelian group, is a polynomial in n for all sufficiently large n . We follow a proof due to Nathanson and Ruzsa as presented in the notes Introduction to Additive Combinatorics by Timothy Gowers for the University of Cambridge.

```

theory Khovanskii
imports
  Complex-Main
  FiniteProduct
  HOL-Library.Equipollence
  Pluenecke-Ruzsa-Inequality.Pluenecke-Ruzsa-Inequality
  Bernoulli.Bernoulli — sums of a fixed power are polynomials
  HOL-Analysis.Weierstrass-Theorems — needed for polynomial function
  HOL-Library.List-Lenlexorder — lexicographic ordering for the type nat

```

```

list

```

```

begin

```

The sum of the elements of a list

```

abbreviation  $\sigma \equiv \text{sum-list}$ 

```

Related to the nsets of Ramsey.thy, but simpler

```

definition finsets :: ['a set, nat]  $\Rightarrow$  'a set set
where finsets A n  $\equiv$  {N. N  $\subseteq$  A  $\wedge$  card N = n}

```

```

lemma card-finsets: finite N  $\implies$  card (finsets N k) = card N choose k
<proof>

```

```

lemma sorted-map-plus-iff [simp]:
fixes a::'a::linordered-cancel-ab-semigroup-add
shows sorted (map ((+) a) xs)  $\longleftrightarrow$  sorted xs
<proof>

```

```

lemma distinct-map-plus-iff [simp]:
fixes a::'a::linordered-cancel-ab-semigroup-add
shows distinct (map ((+) a) xs)  $\longleftrightarrow$  distinct xs
<proof>

```

2.1 Arithmetic operations on lists, pointwise on the elements

Weak type class properties. Cancellation is difficult to arrange because of complications when lists differ in length.

```

instantiation list :: (plus) plus

```

```

begin

```

```

definition plus-list  $\equiv$  map2 (+)

```

```

instance<proof>

```

```

end

```

```

lemma length-plus-list [simp]:
fixes xs :: 'a::plus list
shows length (xs+ys) = min (length xs) (length ys)
<proof>

```

```

lemma plus-Nil [simp]: [] + xs = []
<proof>

```

lemma *plus-Cons*: $(y \# ys) + (x \# xs) = (y+x) \# (ys+xs)$
 ⟨*proof*⟩

lemma *nth-plus-list* [*simp*]:
fixes $xs :: 'a::plus\ list$
assumes $i < length\ xs\ i < length\ ys$
shows $(xs+ys)!i = xs!i + ys!i$
 ⟨*proof*⟩

instantiation $list :: (minus)\ minus$
begin
definition *minus-list* $\equiv map2\ (-)$
instance⟨*proof*⟩
end

lemma *length-minus-list* [*simp*]:
fixes $xs :: 'a::minus\ list$
shows $length\ (xs-ys) = min\ (length\ xs)\ (length\ ys)$
 ⟨*proof*⟩

lemma *minus-Nil* [*simp*]: $[] - xs = []$
 ⟨*proof*⟩

lemma *minus-Cons*: $(y \# ys) - (x \# xs) = (y-x) \# (ys-xs)$
 ⟨*proof*⟩

lemma *nth-minus-list* [*simp*]:
fixes $xs :: 'a::minus\ list$
assumes $i < length\ xs\ i < length\ ys$
shows $(xs-ys)!i = xs!i - ys!i$
 ⟨*proof*⟩

instance $list :: (ab-semigroup-add)\ ab-semigroup-add$
 ⟨*proof*⟩

2.2 The pointwise ordering on two equal-length lists of natural numbers

Gowers uses the usual symbol (\leq) for his pointwise ordering. In our development, \leq is the lexicographic ordering and \sqsubseteq is the pointwise ordering.

definition *pointwise-le* :: $[nat\ list,\ nat\ list] \Rightarrow bool$ (**infixr** \sqsubseteq 50)
where $x \sqsubseteq y \equiv list-all2\ (\leq)\ x\ y$

definition *pointwise-less* :: $[nat\ list,\ nat\ list] \Rightarrow bool$ (**infixr** \triangleleft 50)
where $x \triangleleft y \equiv x \sqsubseteq y \wedge x \neq y$

lemma *pointwise-le-iff-nth*:

$x \sqsubseteq y \longleftrightarrow \text{length } x = \text{length } y \wedge (\forall i < \text{length } x. x!i \leq y!i)$
(proof)

lemma *pointwise-le-iff*:

$x \sqsubseteq y \longleftrightarrow \text{length } x = \text{length } y \wedge (\forall (i,j) \in \text{set } (\text{zip } x \ y). i \leq j)$
(proof)

lemma *pointwise-append-le-iff* [simp]: $u @ x \sqsubseteq u @ y \longleftrightarrow x \sqsubseteq y$
(proof)

lemma *pointwise-le-refl* [iff]: $x \sqsubseteq x$
(proof)

lemma *pointwise-le-antisym*: $\llbracket x \sqsubseteq y; y \sqsubseteq x \rrbracket \Longrightarrow x = y$
(proof)

lemma *pointwise-le-trans*: $\llbracket x \sqsubseteq y; y \sqsubseteq z \rrbracket \Longrightarrow x \sqsubseteq z$
(proof)

lemma *pointwise-le-Nil* [simp]: $\text{Nil} \sqsubseteq x \longleftrightarrow x = \text{Nil}$
(proof)

lemma *pointwise-le-Nil2* [simp]: $x \sqsubseteq \text{Nil} \longleftrightarrow x = \text{Nil}$
(proof)

lemma *pointwise-le-iff-less-equal*: $x \sqsubseteq y \longleftrightarrow x \triangleleft y \vee x = y$
(proof)

lemma *pointwise-less-iff*:

$x \triangleleft y \longleftrightarrow x \sqsubseteq y \wedge (\exists (i,j) \in \text{set } (\text{zip } x \ y). i < j)$
(proof)

lemma *pointwise-less-iff2*: $x \triangleleft y \longleftrightarrow x \sqsubseteq y \wedge (\exists k < \text{length } x. x!k < y ! k)$
(proof)

lemma *pointwise-less-Nil* [simp]: $\neg \text{Nil} \triangleleft x$
(proof)

lemma *pointwise-less-Nil2* [simp]: $\neg x \triangleleft \text{Nil}$
(proof)

lemma *zero-pointwise-le-iff* [simp]: $\text{replicate } r \ 0 \sqsubseteq x \longleftrightarrow \text{length } x = r$
(proof)

lemma *pointwise-le-imp-σ*:

assumes $xs \sqsubseteq ys$ **shows** $\sigma \ xs \leq \sigma \ ys$
(proof)

lemma *sum-list-plus*:

fixes $xs :: 'a::comm-monoid-add list$
assumes $length\ xs = length\ ys$ **shows** $\sigma (xs + ys) = \sigma xs + \sigma ys$
 $\langle proof \rangle$

lemma *sum-list-minus*:

assumes $xs \trianglelefteq ys$ **shows** $\sigma (ys - xs) = \sigma ys - \sigma xs$
 $\langle proof \rangle$

2.3 Pointwise minimum and maximum of a set of lists

definition *min-pointwise* :: $[nat, nat\ list\ set] \Rightarrow nat\ list$

where $min\text{-}pointwise \equiv \lambda r\ U. map\ (\lambda i. Min\ ((\lambda u. u!i) ' U))\ [0..<r]$

lemma *min-pointwise-le*: $\llbracket u \in U; finite\ U \rrbracket \Longrightarrow min\text{-}pointwise\ (length\ u)\ U \trianglelefteq u$
 $\langle proof \rangle$

lemma *min-pointwise-ge-iff*:

assumes $finite\ U\ U \neq \{\}$ $\bigwedge u. u \in U \Longrightarrow length\ u = r$ $length\ x = r$
shows $x \trianglelefteq min\text{-}pointwise\ r\ U \longleftrightarrow (\forall u \in U. x \trianglelefteq u)$
 $\langle proof \rangle$

definition *max-pointwise* :: $[nat, nat\ list\ set] \Rightarrow nat\ list$

where $max\text{-}pointwise \equiv \lambda r\ U. map\ (\lambda i. Max\ ((\lambda u. u!i) ' U))\ [0..<r]$

lemma *max-pointwise-ge*: $\llbracket u \in U; finite\ U \rrbracket \Longrightarrow u \trianglelefteq max\text{-}pointwise\ (length\ u)\ U$
 $\langle proof \rangle$

lemma *max-pointwise-le-iff*:

assumes $finite\ U\ U \neq \{\}$ $\bigwedge u. u \in U \Longrightarrow length\ u = r$ $length\ x = r$
shows $max\text{-}pointwise\ r\ U \trianglelefteq x \longleftrightarrow (\forall u \in U. u \trianglelefteq x)$
 $\langle proof \rangle$

lemma *max-pointwise-mono*:

assumes $X' \subseteq X$ $finite\ X\ X' \neq \{\}$
shows $max\text{-}pointwise\ r\ X' \trianglelefteq max\text{-}pointwise\ r\ X$
 $\langle proof \rangle$

lemma *pointwise-le-plus*: $\llbracket xs \trianglelefteq ys; length\ ys \leq length\ zs \rrbracket \Longrightarrow xs \trianglelefteq ys+zs$
 $\langle proof \rangle$

lemma *pairwise-minus-cancel*: $\llbracket z \trianglelefteq x; z \trianglelefteq y; x - z = y - z \rrbracket \Longrightarrow x = y$
 $\langle proof \rangle$

2.4 A locale to fix the finite subset $A \subseteq G$

locale *Khovanskii* = *additive-abelian-group* +

fixes $A :: 'a\ set$

assumes $AsubG: A \subseteq G$ **and** $finA: finite\ A$

begin

finite products of a group element

definition $Gmult :: 'a \Rightarrow nat \Rightarrow 'a$
where $Gmult\ a\ n \equiv (((\oplus)a) \overset{\sim}{\sim} n)\ \mathbf{0}$

lemma $Gmult-0$ [simp]: $Gmult\ a\ 0 = \mathbf{0}$
 ⟨proof⟩

lemma $Gmult-1$ [simp]: $a \in G \Longrightarrow Gmult\ a\ (Suc\ 0) = a$
 ⟨proof⟩

lemma $Gmult-Suc$ [simp]: $Gmult\ a\ (Suc\ n) = a \oplus Gmult\ a\ n$
 ⟨proof⟩

lemma $Gmult-in-G$ [simp,intro]: $a \in G \Longrightarrow Gmult\ a\ n \in G$
 ⟨proof⟩

lemma $Gmult-add-add$:
assumes $a \in G$
shows $Gmult\ a\ (m+n) = Gmult\ a\ m \oplus Gmult\ a\ n$
 ⟨proof⟩

lemma $Gmult-add-diff$:
assumes $a \in G$
shows $Gmult\ a\ (n+k) \ominus Gmult\ a\ n = Gmult\ a\ k$
 ⟨proof⟩

lemma $Gmult-diff$:
assumes $a \in G\ n \leq m$
shows $Gmult\ a\ m \ominus Gmult\ a\ n = Gmult\ a\ (m-n)$
 ⟨proof⟩

Mapping elements of A to their numeric subscript

abbreviation $idx \equiv to-nat-on\ A$

The elements of A in order

definition $aA :: 'a\ list$
where $aA \equiv map\ (from-nat-into\ A)\ [0..<card\ A]$

definition $\alpha :: nat\ list \Rightarrow 'a$
where $\alpha \equiv \lambda x. fincomp\ (\lambda i. Gmult\ (aA!i)\ (x!i))\ \{..<card\ A\}$

The underlying assumption is $length\ y = length\ x$

definition $useless :: nat\ list \Rightarrow bool$
where $useless\ x \equiv \exists y < x. \sigma\ y = \sigma\ x \wedge \alpha\ y = \alpha\ x \wedge length\ y = length\ x$

abbreviation $useful\ x \equiv \neg\ useless\ x$

lemma $alpha-replicate-0$ [simp]: $\alpha\ (replicate\ (card\ A)\ 0) = \mathbf{0}$
 ⟨proof⟩

lemma *idx-less-cardA*:

assumes $a \in A$ **shows** $\text{idx } a < \text{card } A$
<proof>

lemma *aA-idx-eq* [simp]:

assumes $a \in A$ **shows** $aA ! (\text{idx } a) = a$
<proof>

lemma *set-aA*: $\text{set } aA = A$

<proof>

lemma *nth-aA-in-G* [simp]: $i < \text{card } A \implies aA!i \in G$

<proof>

lemma *alpha-in-G* [iff]: $\alpha x \in G$

<proof>

lemma *Gmult-in-PiG* [simp]: $(\lambda i. \text{Gmult } (aA!i) (f i)) \in \{..<\text{card } A\} \rightarrow G$

<proof>

lemma *alpha-plus*:

assumes $\text{length } x = \text{card } A$ $\text{length } y = \text{card } A$

shows $\alpha (x + y) = \alpha x \oplus \alpha y$

<proof>

lemma *alpha-minus*:

assumes $y \triangleleft x$ $\text{length } y = \text{card } A$

shows $\alpha (x - y) = \alpha x \ominus \alpha y$

<proof>

2.5 Adding one to a list element

definition *list-incr* :: $\text{nat} \Rightarrow \text{nat list} \Rightarrow \text{nat list}$

where $\text{list-incr } i x \equiv x[i := \text{Suc } (x!i)]$

lemma *list-incr-Nil* [simp]: $\text{list-incr } i [] = []$

<proof>

lemma *list-incr-Cons* [simp]: $\text{list-incr } (Suc i) (k\#ks) = k \# \text{list-incr } i ks$

<proof>

lemma *sum-list-incr* [simp]: $i < \text{length } x \implies \sigma (\text{list-incr } i x) = \text{Suc } (\sigma x)$

<proof>

lemma *length-list-incr* [simp]: $\text{length } (\text{list-incr } i x) = \text{length } x$

<proof>

lemma *nth-le-list-incr*: $i < \text{card } A \implies x!i \leq \text{list-incr } (idx a) x!i$

<proof>

lemma *list-incr-nth-diff*: $i < \text{length } x \implies \text{list-incr } j \ x!i - x!i = (\text{if } i = j \text{ then } 1 \text{ else } 0)$

<proof>

2.6 The set of all r -tuples that sum to n

definition *length-sum-set* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat list set}$

where $\text{length-sum-set } r \ n \equiv \{x. \text{length } x = r \wedge \sigma \ x = n\}$

lemma *length-sum-set-Nil* [*simp*]: $\text{length-sum-set } 0 \ n = (\text{if } n=0 \text{ then } \{\}\ \text{else } \{\})$

<proof>

lemma *length-sum-set-Suc* [*simp*]: $k \# \text{ks} \in \text{length-sum-set } (\text{Suc } r) \ n \longleftrightarrow (\exists m. \text{ks} \in \text{length-sum-set } r \ m \wedge n = m+k)$

<proof>

lemma *length-sum-set-Suc-egpoll*: $\text{length-sum-set } (\text{Suc } r) \ n \approx \text{Sigma } \{..n\} \ (\lambda i. \text{length-sum-set } r \ (n-i))$ (**is** $?L \approx ?R$)

<proof>

lemma *finite-length-sum-set*: $\text{finite } (\text{length-sum-set } r \ n)$

<proof>

lemma *card-length-sum-set*: $\text{card } (\text{length-sum-set } (\text{Suc } r) \ n) = (\sum i \leq n. \text{card } (\text{length-sum-set } r \ (n-i)))$

<proof>

lemma *sum-up-index-split'*:

assumes $N \leq n$ **shows** $(\sum i \leq n. f \ i) = (\sum i \leq n-N. f \ i) + (\sum i = \text{Suc } (n-N)..n. f \ i)$

<proof>

lemma *sum-invert*: $N \leq n \implies (\sum i = \text{Suc } (n - N)..n. f \ (n - i)) = (\sum j < N. f \ j)$

<proof>

lemma *sum-diff-split*:

assumes $N \leq n$

shows $(\sum i \leq n - N. \text{real } (n - i) \wedge j) = (\sum i \leq n. \text{real } i \wedge j) - (\sum i < N. \text{real } i \wedge j)$

<proof>

lemma *real-polynomial-function-length-sum-set*:

$\exists p. \text{real-polynomial-function } p \wedge (\forall n > 0. \text{real } (\text{card } (\text{length-sum-set } r \ n)) = p \ (\text{real } n))$

<proof>

lemma *all-zeroes-replicate*: $\text{length-sum-set } r \ 0 = \{\text{replicate } r \ 0\}$
 ⟨proof⟩

lemma *length-sum-set-Suc-eq-UN*: $\text{length-sum-set } r \ (\text{Suc } n) = (\bigcup i < r. \text{list-incr } i \text{ ` length-sum-set } r \ n)$
 ⟨proof⟩

lemma *alpha-list-incr*:
assumes $a \in A \ x \in \text{length-sum-set } (\text{card } A) \ n$
shows $\alpha \ (\text{list-incr } (\text{id } a) \ x) = a \oplus \alpha \ x$
 ⟨proof⟩

lemma *sumset-iterated-enum*:
defines $r \equiv \text{card } A$
shows $\text{sumset-iterated } A \ n = \alpha \ \text{` length-sum-set } r \ n$
 ⟨proof⟩

2.7 Lemma 2.7 in Gowers's notes

The following lemma corresponds to a key fact about the cardinality of the set of all sums of n many elements of A , stated before Gowers's Lemma 2.7.

lemma *card-sumset-iterated-length-sum-set-useful*:
defines $r \equiv \text{card } A$
shows $\text{card}(\text{sumset-iterated } A \ n) = \text{card}(\text{length-sum-set } r \ n \cap \{x. \text{useful } x\})$
 (is $\text{card } ?L = \text{card } ?R$)
 ⟨proof⟩

The following lemma corresponds to Lemma 2.7 in Gowers's notes.

lemma *useless-leq-useless*:
defines $r \equiv \text{card } A$
assumes *useless* x **and** $x \leq y$ **and** $\text{length } x = r$
shows *useless* y
 ⟨proof⟩

inductive-set *minimal-elements for* U
where $\llbracket x \in U; \bigwedge y. y \in U \implies \neg y \triangleleft x \rrbracket \implies x \in \text{minimal-elements } U$

lemma *pointwise-less-imp-σ*:
assumes $xs \triangleleft ys$ **shows** $\sigma \ xs < \sigma \ ys$
 ⟨proof⟩

lemma *wf-measure-σ*: *wf* (*inv-image less-than* σ)
 ⟨proof⟩

lemma *WFP*: *wfP* (\triangleleft)
 ⟨proof⟩

The following is a direct corollary of the above lemma, i.e. a corollary of Lemma 2.7 in Gowers's notes.

corollary *useless-iff*:

assumes $\text{length } x = \text{card } A$

shows $\text{useless } x \iff (\exists x' \in \text{minimal-elements } (\text{Collect useless}). x' \trianglelefteq x)$ (is -=?R)

<proof>

2.8 The set of minimal elements of a set of r -tuples is finite

The following general finiteness claim corresponds to Lemma 2.8 in Gowers's notes and is key to the main proof.

lemma *minimal-elements-set-tuples-finite*:

assumes $Ur: \bigwedge x. x \in U \implies \text{length } x = r$

shows *finite* (*minimal-elements* U)

<proof>

2.9 Towards Lemma 2.9 in Gowers's notes

Increasing sequences

fun *augmentum* :: *nat list* \Rightarrow *nat list*

where *augmentum* [] = []

| *augmentum* ($n\#ns$) = $n \# \text{map } ((+)n) (\text{augmentum } ns)$

definition *dementum*:: *nat list* \Rightarrow *nat list*

where *dementum* $xs \equiv xs - (0\#xs)$

lemma *dementum-Nil* [*simp*]: *dementum* [] = []

<proof>

lemma *zero-notin-augmentum* [*simp*]: $0 \notin \text{set } ns \implies 0 \notin \text{set } (\text{augmentum } ns)$

<proof>

lemma *length-augmentum* [*simp*]: $\text{length } (\text{augmentum } xs) = \text{length } xs$

<proof>

lemma *sorted-augmentum* [*simp*]: $0 \notin \text{set } ns \implies \text{sorted } (\text{augmentum } ns)$

<proof>

lemma *distinct-augmentum* [*simp*]: $0 \notin \text{set } ns \implies \text{distinct } (\text{augmentum } ns)$

<proof>

lemma *augmentum-subset-sum-list*: $\text{set } (\text{augmentum } ns) \subseteq \{..\sigma ns\}$

<proof>

lemma *sum-list-augmentum*: $\sigma ns \in \text{set } (\text{augmentum } ns) \iff \text{length } ns > 0$

<proof>

lemma *length-dementum* [*simp*]: $\text{length} (\text{dementum } xs) = \text{length } xs$
 ⟨*proof*⟩

lemma *sorted-imp-pointwise*:
assumes *sorted* ($xs@[n]$)
shows $0 \# xs \leq xs @ [n]$
 ⟨*proof*⟩

lemma *sum-list-dementum*:
assumes *sorted* ($xs@[n]$)
shows $\sigma (\text{dementum } (xs@[n])) = n$
 ⟨*proof*⟩

lemma *augmentum-cancel*: $\text{map } ((+)k) (\text{augmentum } ns) - (k \# \text{map } ((+)k) (\text{augmentum } ns)) = ns$
 ⟨*proof*⟩

lemma *dementum-augmentum* [*simp*]:
assumes $0 \notin \text{set } ns$
shows $(\text{dementum} \circ \text{sorted-list-of-set}) ((\text{set} \circ \text{augmentum}) ns) = ns$ (**is** ?*L* *ns* = -)
 ⟨*proof*⟩

lemma *dementum-nonzero*:
assumes *ns: sorted-wrt* ($<$) *ns* **and** $0 \notin \text{set } ns$
shows $0 \notin \text{set} (\text{dementum } ns)$
 ⟨*proof*⟩

lemma *nth-augmentum* [*simp*]: $i < \text{length } ns \implies \text{augmentum } ns!i = (\sum_{j \leq i} ns!j)$
 ⟨*proof*⟩

lemma *augmentum-dementum* [*simp*]:
assumes $0 \notin \text{set } ns$ *sorted* *ns*
shows $\text{augmentum} (\text{dementum } ns) = ns$
 ⟨*proof*⟩

The following lemma corresponds to Lemma 2.9 in Gowers’s notes. The proof involves introducing bijective maps between r -tuples that fulfill certain properties/ r -tuples and subsets of naturals, so as to show the cardinality claim.

lemma *bound-sum-list-card*:
assumes $r > 0$ **and** $n \geq \sigma x'$ **and** *len* x' : $\text{length } x' = r$
defines $S \equiv \{x. x' \leq x \wedge \sigma x = n\}$
shows $\text{card } S = (n - \sigma x' + r - 1)$ *choose* $(r-1)$
 ⟨*proof*⟩

2.10 Towards the main theorem

lemma *extend-tuple*:

assumes $\sigma xs \leq n$ $length\ xs \neq 0$
obtains ys **where** $\sigma ys = n$ $xs \trianglelefteq ys$
 ⟨proof⟩

lemma *extend-preserving*:
assumes $\sigma xs \leq n$ $length\ xs > 1$ $i < length\ xs$
obtains ys **where** $\sigma ys = n$ $xs \trianglelefteq ys$ $ys!i = xs!i$
 ⟨proof⟩

The proof of the main theorem will make use of the inclusion-exclusion formula, in addition to the previously shown results.

theorem *Khovanskii*:
assumes $card\ A > 1$
defines $f \equiv \lambda n. card(sumset-iterated\ A\ n)$
obtains $N\ p$ **where** *real-polynomial-function* $p \wedge n. n \geq N \implies real\ (f\ n) = p$
 (*real* n)
 ⟨proof⟩

end

end

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