

# Deriving generic class instances for datatypes

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## Abstract

We provide a framework for automatically deriving instances for generic type classes. Our approach is inspired by Haskell’s *generic-deriving* package [1] and Scala’s *shapeless* library [2].

In addition to generating the code for type class functions, we also attempt to automatically prove type class laws for these instances. As of now, however, some manual proofs are still required for recursive datatypes.

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## 1 Tagged Sum-of-Products Representation

This theory sets up a version of the sum-of-products representation that includes constructor and selector names. For an example of a type class that uses this representation see `Derive_Show`.

```
theory Tagged-Prod-Sum  
imports Main  
begin
```

**context begin**

**qualified datatype** ('a, 'b) prod = Prod string option string option 'a 'b  
**qualified datatype** ('a, 'b) sum = Inl string option 'a | Inr string option 'b

**qualified definition** fst **where** fst p = (case p of (Prod - - a -) ⇒ a)  
**qualified definition** snd **where** snd p = (case p of (Prod - - - b) ⇒ b)  
**qualified definition** sel-name-fst **where** sel-name-fst p = (case p of (Prod s - -) ⇒ s)  
**qualified definition** sel-name-snd **where** sel-name-snd p = (case p of (Prod - s -) ⇒ s)

**qualified definition** constr-name **where** constr-name x = (case x of (Inl s -) ⇒ s | (Inr s -) ⇒ s)

**end**

**lemma** measure-tagged-fst[measure-function]: is-measure f ⇒ is-measure (λ p. f (Tagged-Prod-Sum.fst p))  
⟨proof⟩

**lemma** measure-tagged-snd[measure-function]: is-measure f ⇒ is-measure (λ p. f (Tagged-Prod-Sum.snd p))  
⟨proof⟩

**lemma** size-tagged-prod-simp:  
Tagged-Prod-Sum.prod.size-prod f g p = f (Tagged-Prod-Sum.fst p) + g (Tagged-Prod-Sum.snd p) + Suc 0  
⟨proof⟩

**lemma** size-tagged-sum-simp:  
Tagged-Prod-Sum.sum.size-sum f g x = (case x of Tagged-Prod-Sum.Inl - a ⇒ f a + Suc 0 | Tagged-Prod-Sum.Inr - b ⇒ g b + Suc 0)  
⟨proof⟩

**lemma** size-tagged-prod-measure:  
is-measure f ⇒ is-measure g ⇒ is-measure (Tagged-Prod-Sum.prod.size-prod f g)  
⟨proof⟩

**lemma** size-tagged-sum-measure:  
is-measure f ⇒ is-measure g ⇒ is-measure (Tagged-Prod-Sum.sum.size-sum

```
f g)
⟨proof⟩
```

```
end
```

## 2 Derive

This theory includes the Isabelle/ML code needed for the derivation and exports the two keywords `derive_generic` and `derive_generic_setup`.

```
theory Derive
```

```
  imports Main Tagged-Prod-Sum
```

```
  keywords derive-generic derive-generic-setup :: thy-goal
```

```
begin
```

```
context begin
```

```
qualified definition iso :: ('a ⇒ 'b) ⇒ ('b ⇒ 'a) ⇒ bool where  
iso from to = ((∀ a. to (from a) = a) ∧ (∀ b. from (to b) = b))
```

```
lemma iso-intro: (∧a. to (from a) = a) ⇒ (∧b. from (to b) = b) ⇒ iso  
from to  
  ⟨proof⟩
```

```
end
```

```
⟨ML⟩
```

```
end
```

## 3 Examples

### 3.1 Example Datatypes

```
theory Derive-Datatypes
```

```
  imports Main
```

```
begin
```

```
datatype simple = A (num: nat) | B (left:nat) (right:nat) | C
```

```
datatype ('a,'b) either = L 'a | R 'b
```

```
datatype 'a tree = Leaf | Node 'a 'a tree 'a tree
```

```
datatype even-nat = Even-Zero | Even-Succ odd-nat  
and odd-nat = Odd-Succ even-nat
```

```
datatype ('a,'b) exp = Term ('a,'b) trm | Sum (left:('a,'b) trm) (right:('a,'b)  
exp)  
and ('a,'b) trm = Factor ('a,'b) fct | Prod ('a,'b) fct ('a,'b) trm  
and ('a,'b) fct = Const 'a | Var (v:'b) | Expr ('a,'b) exp  
  
end
```

### 3.2 Equality

```
theory Derive-Eq  
imports Main ../Derive Derive-Datatypes  
begin
```

```
class eq =  
fixes eq :: 'a  $\Rightarrow$  'a  $\Rightarrow$  bool
```

```
instantiation nat and unit:: eq  
begin  
definition eq-nat : eq (x::nat) y  $\longleftrightarrow$  x = y  
definition eq-unit-def: eq (x::unit) y  $\longleftrightarrow$  True  
instance <proof>  
end
```

```
instantiation prod and sum :: (eq, eq) eq  
begin  
definition eq-prod-def: eq x y  $\longleftrightarrow$  (eq (fst x) (fst y))  $\wedge$  (eq (snd x) (snd  
y))  
definition eq-sum-def: eq x y = (case x of Inl a  $\Rightarrow$  (case y of Inl b  $\Rightarrow$  eq  
a b | Inr b  $\Rightarrow$  False)  
| Inr a  $\Rightarrow$  (case y of Inl b  $\Rightarrow$  False | Inr  
b  $\Rightarrow$  eq a b))  
  
instance <proof>  
end
```

**derive-generic** *eq simple*  $\langle proof \rangle$

**lemma** *eq (A 4) (A 4)*  $\langle proof \rangle$

**lemma** *eq (A 6) (A 4)*  $\longleftrightarrow False \langle proof \rangle$

**lemma** *eq C C*  $\langle proof \rangle$

**lemma** *eq (B 4 5) (B 4 5)*  $\langle proof \rangle$

**lemma** *eq (B 4 4) (A 3)*  $\longleftrightarrow False \langle proof \rangle$

**lemma** *eq C (A 4)*  $\longleftrightarrow False \langle proof \rangle$

**derive-generic** *eq either*  $\langle proof \rangle$

**lemma** *eq (L (3::nat)) (R 3)*  $\longleftrightarrow False \langle proof \rangle$

**lemma** *eq (L (3::nat)) (L 3)*  $\langle proof \rangle$

**lemma** *eq (L (3::nat)) (L 4)*  $\longleftrightarrow False \langle proof \rangle$

**derive-generic** *eq list*  $\langle proof \rangle$

**lemma** *eq ([ ]::(nat list)) [ ]*  $\langle proof \rangle$

**lemma** *eq ([1,2,3]::(nat list)) [1,2,3]*  $\langle proof \rangle$

**lemma** *eq [(1::nat)] [1,2]*  $\longleftrightarrow False \langle proof \rangle$

**derive-generic** *eq tree*  $\langle proof \rangle$

**lemma** *eq Leaf Leaf*  $\langle proof \rangle$

**lemma** *eq (Node (1::nat) Leaf Leaf) Leaf*  $\longleftrightarrow False \langle proof \rangle$

**lemma** *eq (Node (1::nat) Leaf Leaf) (Node (1::nat) Leaf Leaf)*  $\langle proof \rangle$

**lemma** *eq (Node (1::nat) (Node 2 Leaf Leaf) (Node 3 Leaf Leaf)) (Node (1::nat) (Node 2 Leaf Leaf) (Node 4 Leaf Leaf))*  
 $\longleftrightarrow False \langle proof \rangle$

**derive-generic** *eq even-nat*  $\langle proof \rangle$

**derive-generic** *eq exp*  $\langle proof \rangle$

**lemma** *eq Even-Zero Even-Zero*  $\langle proof \rangle$

**lemma** *eq Even-Zero (Even-Succ (Odd-Succ Even-Zero))*  $\longleftrightarrow False \langle proof \rangle$

```

lemma eq (Odd-Succ (Even-Succ (Odd-Succ Even-Zero))) (Odd-Succ (Even-Succ
(Odd-Succ Even-Zero))) ⟨proof⟩
lemma eq (Odd-Succ (Even-Succ (Odd-Succ Even-Zero))) (Odd-Succ (Even-Succ
(Odd-Succ (Even-Succ (Odd-Succ Even-Zero))))))
  ⟷ False ⟨proof⟩

lemma eq (Const (1::nat)) (Const (1::nat)) ⟨proof⟩
lemma eq (Const (1::nat)) (Var (1::nat)) ⟷ False ⟨proof⟩
lemma eq (Term (Prod (Const (1::nat)) (Factor (Const (2::nat))))) (Term
(Prod (Const (1::nat)) (Factor (Const (2::nat)))))
  ⟨proof⟩
lemma eq (Term (Prod (Const (1::nat)) (Factor (Const (2::nat))))) (Term
(Prod (Const (1::nat)) (Factor (Const (3::nat)))))
  ⟷ False ⟨proof⟩

```

**end**

### 3.3 Encoding

```

theory Derive-Encode
imports Main ../Derive Derive-Datatypes
begin

class encodeable =
  fixes encode :: 'a ⇒ bool list

instantiation nat and unit:: encodeable
begin
  fun encode-nat :: nat ⇒ bool list where
    encode-nat 0 = [] |
    encode-nat (Suc n) = True # (encode n)

  definition encode-unit: encode (x::unit) = []
  instance ⟨proof⟩
end

instantiation prod and sum :: (encodeable, encodeable) encodeable
begin
  definition encode-prod-def: encode x = append (encode (fst x)) (encode
(snd x))
  definition encode-sum-def: encode x = (case x of Inl a ⇒ False # encode
a
  | Inr a ⇒ True # encode a)

```

```

instance ⟨proof⟩
end

derive-generic encodeable simple ⟨proof⟩
derive-generic encodeable either ⟨proof⟩

lemma encode (B 3 4) = [True, False, True, True, True, True, True, True, True, True] ⟨proof⟩
lemma encode C = [True, True] ⟨proof⟩
lemma encode (R (3::nat)) = [True, True, True, True] ⟨proof⟩

derive-generic encodeable list ⟨proof⟩
derive-generic encodeable tree ⟨proof⟩

lemma encode [1,2,3,4::nat]
  = [True, True, True, True, True, True, True, True, True, True, True, True, True, True, True, False] ⟨proof⟩
lemma encode (Node (3::nat) (Node 1 Leaf Leaf) (Node 2 Leaf Leaf))
  = [True, True, True, True, True, True, False, False, True, True, True, False, False] ⟨proof⟩

derive-generic encodeable even-nat ⟨proof⟩
derive-generic encodeable exp ⟨proof⟩

lemma encode (Odd-Succ (Even-Succ (Odd-Succ Even-Zero)))
  = [True, False, True, True, False, False] ⟨proof⟩
lemma encode (Term (Prod (Const (1::nat)) (Factor (Const (2::nat)))))
  = [False, False, True, False, True, True, True, False, True, True, False, False, True, True, True, False, True, True]
  ⟨proof⟩

end

3.4 Algebraic Classes

theory Derive-Algebra
imports Main ../Derive Derive-Datatypes
begin

class semigroup =

```

```

fixes mult :: 'a ⇒ 'a ⇒ 'a (infixl ⊗ 70)

class monoidl = semigroup +
fixes neutral :: 'a (1)

class group = monoidl +
  fixes inverse :: 'a ⇒ 'a

instantiation nat and unit:: semigroup
begin
  definition mult-nat : mult (x::nat) y = x + y
  definition mult-unit-def: mult (x::unit) y = x
  instance ⟨proof⟩
end
instantiation nat and unit:: monoidl
begin
  definition neutral-nat : neutral = (0::nat)
  definition neutral-unit-def: neutral = ()
  instance ⟨proof⟩
end

instantiation nat and unit:: group
begin
  definition inverse-nat : inverse (i::nat) = 1 - i
  definition inverse-unit-def: inverse u = ()
  instance ⟨proof⟩
end

instantiation prod and sum :: (semigroup, semigroup) semigroup
begin
  definition mult-prod-def: x ⊗ y = (fst x ⊗ fst y, snd x ⊗ snd y)
  definition mult-sum-def: x ⊗ y = (case x of Inl a ⇒ (case y of Inl b ⇒
Inl (a ⊗ b) | Inr b ⇒ Inl a)
  | Inr a ⇒ (case y of Inl b ⇒ Inr a | Inr
b ⇒ Inr (a ⊗ b)))
  instance ⟨proof⟩
end

instantiation prod and sum :: (monoidl, monoidl) monoidl
begin

```



```

definition neutral-prod-def: neutral = (neutral,neutral)
definition neutral-sum-def: neutral = Inl neutral
instance  $\langle$ proof $\rangle$ 
end

instantiation prod and sum :: (group, group) group
begin
  definition inverse-prod-def: inverse p = (inverse (fst p), inverse (snd p))
  definition inverse-sum-def: inverse x = (case x of Inl a  $\Rightarrow$  (Inl (inverse a))
  | Inr b  $\Rightarrow$  Inr (inverse b))

  instance  $\langle$ proof $\rangle$ 
end

```

```

derive-generic semigroup simple  $\langle$ proof $\rangle$ 
derive-generic monoidl simple  $\langle$ proof $\rangle$ 
derive-generic group simple  $\langle$ proof $\rangle$ 

```

```

lemma (B 1 6)  $\otimes$  (B 4 5) = B 4 11  $\langle$ proof $\rangle$ 
lemma (A 2)  $\otimes$  (A 3) = A 5  $\langle$ proof $\rangle$ 
lemma (B 1 6)  $\otimes$  1 = B 0 6  $\langle$ proof $\rangle$ 

```

```

derive-generic group either  $\langle$ proof $\rangle$ 

```

```

lemma (L 3)  $\otimes$  ((L 4)::(nat,nat) either) = L 7  $\langle$ proof $\rangle$ 
lemma (R (2::nat))  $\otimes$  (L (3::nat)) = R 2  $\langle$ proof $\rangle$ 

```

```

derive-generic semigroup list  $\langle$ proof $\rangle$ 
derive-generic monoidl list  $\langle$ proof $\rangle$ 
derive-generic group list  $\langle$ proof $\rangle$ 
derive-generic semigroup tree  $\langle$ proof $\rangle$ 
derive-generic monoidl tree  $\langle$ proof $\rangle$ 
derive-generic group tree  $\langle$ proof $\rangle$ 

```

```

lemma [1,2,3,4::nat]  $\otimes$  [1,2,3] = [2,4,6,4]  $\langle$ proof $\rangle$ 
lemma inverse [1,2,3::nat] = [0,0,0]  $\langle$ proof $\rangle$ 

```

```

derive-generic semigroup even-nat ⟨proof⟩
derive-generic monoidl even-nat ⟨proof⟩
derive-generic group even-nat ⟨proof⟩
derive-generic semigroup exp ⟨proof⟩

```

```

instantiation exp and trm and fct :: (monoidl,monoidl) monoidl
begin
  definition neutral-fct where neutral-fct = Const neutral
  definition neutral-trm where neutral-trm = Factor neutral
  definition neutral-exp where neutral-exp = Term neutral
  instance ⟨proof⟩
end

```

⟨ML⟩

```

derive-generic group exp ⟨proof⟩

```

```

lemma (Odd-Succ (Even-Succ (Odd-Succ Even-Zero))) ⊗ (Odd-Succ Even-Zero)

```

```

  = Odd-Succ (Even-Succ (Odd-Succ Even-Zero)) ⟨proof⟩

```

```

lemma inverse (Odd-Succ Even-Zero) = Odd-Succ Even-Zero ⟨proof⟩

```

```

lemma (Term (Prod ((Const 1)::(nat, nat) fct) (Factor (Const (2::nat)))))

```

```

  ⊗ (Term (Prod (Const (2::nat)) (Factor ((Const 2)::(nat, nat) fct))))

```

```

  = Term (Prod (Const 3) (Factor (Const 4))) ⟨proof⟩

```

```

end

```

### 3.5 Show

```

theory Derive-Show

```

```

imports Main ../Derive Derive-Datatypes

```

```

begin

```

```

class showable =

```

```

  fixes print :: 'a ⇒ string

```

```

fun string-of-nat :: nat ⇒ string

```

```

where

```

```

  string-of-nat n = (if n < 10 then [(char-of :: nat ⇒ char) (48 + n)] else

```

*string-of-nat* (n div 10) @ [(char-of :: nat ⇒ char) (48 + (n mod 10))]

**instantiation** *nat and unit:: showable*

**begin**

**definition** *print-nat*: print (n::nat) = *string-of-nat* n

**definition** *print-unit*: print (x::unit) = ""

**instance** ⟨*proof*⟩

**end**

**instantiation** *Tagged-Prod-Sum.prod and Tagged-Prod-Sum.sum :: (showable, showable) showable*

**begin**

**definition** *print-prod-def*:

print (x::('a,'b) *Tagged-Prod-Sum.prod*) =

(case *Tagged-Prod-Sum.sel-name-fst* x of

None ⇒ (print (*Tagged-Prod-Sum.fst* x))

| Some s ⇒ "(" @ s @ ": " @ (print (*Tagged-Prod-Sum.fst* x)) @ ")")

@

" "

@

(case *Tagged-Prod-Sum.sel-name-snd* x of

None ⇒ (print (*Tagged-Prod-Sum.snd* x))

| Some s ⇒ "(" @ s @ ": " @ (print (*Tagged-Prod-Sum.snd* x)) @ ")")

**definition** *print-sum-def*: print (x::('a,'b) *Tagged-Prod-Sum.sum*) =

(case x of (*Tagged-Prod-Sum.Inl* s a) ⇒ (case s of None ⇒ print a | Some c ⇒ "(" @ c @ " " @ (print a) @ ")")

| (*Tagged-Prod-Sum.Inr* s b) ⇒ (case s of None ⇒ print b | Some c ⇒ "(" @ c @ " " @ (print b) @ ")")

**instance** ⟨*proof*⟩

**end**

**declare** [[*ML-print-depth=30*]]

**derive-generic** (*metadata*) *showable simple* ⟨*proof*⟩

**derive-generic** (*metadata*) *showable either* ⟨*proof*⟩

**value** *print* (A 3)

**value** *print* (B 1 2)

**value** [*simp*] *print* (L (2::nat))

**value** *print* C

```

derive-generic (metadata) showable list ⟨proof⟩
derive-generic (metadata) showable tree ⟨proof⟩

value print [1,2::nat]
value print (Node (3::nat) (Node 1 Leaf Leaf) (Node 2 Leaf Leaf))

derive-generic (metadata) showable even-nat ⟨proof⟩
derive-generic (metadata) showable exp ⟨proof⟩

value print (Odd-Succ (Even-Succ (Odd-Succ Even-Zero)))
value [simp] print (Sum (Factor (Const (0::nat))) (Term (Prod (Const
(1::nat)) (Factor (Const (2::nat)))))))

end

```

## 3.6 Classes with Laws

### 3.6.1 Equality

```

theory Derive-Eq-Laws
  imports Main ../Derive Derive-Datatypes
begin

class eq =
  fixes eq :: 'a ⇒ 'a ⇒ bool
  assumes refl: eq x x and
    sym: eq x y ⇒ eq y x and
    trans: eq x y ⇒ eq y z ⇒ eq x z

derive-generic-setup eq
  ⟨proof⟩

lemma eq-law-eq: eq-class-law eq
  ⟨proof⟩

instantiation nat and unit :: eq
begin
  definition eq-nat-def : eq (x::nat) y ↔ x = y

```

**definition** *eq-unit-def*:  $eq (x::unit) y \longleftrightarrow True$   
**instance**  $\langle proof \rangle$   
**end**

**instantiation** *prod and sum* :: (eq, eq) eq  
**begin**

**definition** *eq-prod-def*:  $eq x y \longleftrightarrow (eq (fst x) (fst y)) \wedge (eq (snd x) (snd y))$

**definition** *eq-sum-def*:  $eq x y = (case\ x\ of\ Inl\ a \Rightarrow (case\ y\ of\ Inl\ b \Rightarrow eq\ a\ b\ | Inr\ b \Rightarrow False) \mid Inr\ a \Rightarrow (case\ y\ of\ Inl\ b \Rightarrow False\ | Inr\ b \Rightarrow eq\ a\ b))$

**instance**  $\langle proof \rangle$   
**end**

**derive-generic** *eq simple*  $\langle proof \rangle$

**lemma**  $eq (A\ 4) (A\ 4) \langle proof \rangle$   
**lemma**  $eq (A\ 6) (A\ 4) \longleftrightarrow False \langle proof \rangle$   
**lemma**  $eq\ C\ C \langle proof \rangle$   
**lemma**  $eq (B\ 4\ 5) (B\ 4\ 5) \langle proof \rangle$   
**lemma**  $eq (B\ 4\ 4) (A\ 3) \longleftrightarrow False \langle proof \rangle$   
**lemma**  $eq\ C (A\ 4) \longleftrightarrow False \langle proof \rangle$

**derive-generic** *eq either*  $\langle proof \rangle$

**lemma**  $eq (L (3::nat)) (R\ 3) \longleftrightarrow False \langle proof \rangle$   
**lemma**  $eq (L (3::nat)) (L\ 3) \langle proof \rangle$   
**lemma**  $eq (L (3::nat)) (L\ 4) \longleftrightarrow False \langle proof \rangle$

**derive-generic** *eq list*  
 $\langle proof \rangle$

**lemma**  $eq ([ ]::(nat\ list)) [ ] \langle proof \rangle$   
**lemma**  $eq ([1,2,3]::(nat\ list)) [1,2,3] \langle proof \rangle$   
**lemma**  $eq [(1::nat)] [1,2] \longleftrightarrow False \langle proof \rangle$

**derive-generic** *eq tree*  
 $\langle proof \rangle$

```

lemma eq Leaf Leaf ⟨proof⟩
lemma eq (Node (1::nat) Leaf Leaf) Leaf  $\longleftrightarrow$  False ⟨proof⟩
lemma eq (Node (1::nat) Leaf Leaf) (Node (1::nat) Leaf Leaf) ⟨proof⟩
lemma eq (Node (1::nat) (Node 2 Leaf Leaf) (Node 3 Leaf Leaf)) (Node
(1::nat) (Node 2 Leaf Leaf) (Node 4 Leaf Leaf))
 $\longleftrightarrow$  False ⟨proof⟩
end

```

### 3.6.2 Algebraic Classes

```

theory Derive-Algebra-Laws
  imports Main ../Derive Derive-Datatypes
begin

```

```

datatype simple-int = A int | B int int | C

```

```

class semigroup =
  fixes mult :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixl  $\otimes$  70)
  assumes assoc: (x  $\otimes$  y)  $\otimes$  z = x  $\otimes$  (y  $\otimes$  z)

```

```

class monoidl = semigroup +
  fixes neutral :: 'a (1)
  assumes neutl : 1  $\otimes$  x = x

```

```

class group = monoidl +
  fixes inverse :: 'a  $\Rightarrow$  'a
  assumes invl: (inverse x)  $\otimes$  x = 1

```

```

definition semigroup-law :: ('a  $\Rightarrow$  'a  $\Rightarrow$  'a)  $\Rightarrow$  bool where
semigroup-law MULT = ( $\forall$  x y z. MULT (MULT x y) z = MULT x (MULT
y z))

```

```

definition monoidl-law :: 'a  $\Rightarrow$  ('a  $\Rightarrow$  'a  $\Rightarrow$  'a)  $\Rightarrow$  bool where
monoidl-law NEUTRAL MULT = (( $\forall$  x. MULT NEUTRAL x = x)  $\wedge$  semi-
group-law MULT)

```

```

definition group-law :: ('a  $\Rightarrow$  'a)  $\Rightarrow$  'a  $\Rightarrow$  ('a  $\Rightarrow$  'a  $\Rightarrow$  'a)  $\Rightarrow$  bool where
group-law INVERSE NEUTRAL MULT = (( $\forall$  x. MULT (INVERSE x) x
= NEUTRAL)  $\wedge$  monoidl-law NEUTRAL MULT)

```

```

lemma transfer-semigroup:
  assumes Derive.iso f g
  shows semigroup-law MULT  $\implies$  semigroup-law ( $\lambda$ x y. g (MULT (f x) (f
y)))
  ⟨proof⟩

```

**lemma** *transfer-monoidl*:  
**assumes** *Derive.iso f g*  
**shows** *monoidl-law NEUTRAL MULT  $\implies$  monoidl-law (g NEUTRAL)*  
 $(\lambda x y. g (MULT (f x) (f y)))$   
 $\langle proof \rangle$

**lemma** *transfer-group*:  
**assumes** *Derive.iso f g*  
**shows** *group-law INVERSE NEUTRAL MULT  $\implies$  group-law  $(\lambda x. g$*   
 $(INVERSE (f x))) (g NEUTRAL) (\lambda x y. g (MULT (f x) (f y)))$   
 $\langle proof \rangle$

**lemma** *semigroup-law-semigroup: semigroup-law mult*  
 $\langle proof \rangle$

**lemma** *monoidl-law-monoidl: monoidl-law neutral mult*  
 $\langle proof \rangle$

**lemma** *group-law-group: group-law inverse neutral mult*  
 $\langle proof \rangle$

**derive-generic-setup** *semigroup*  
 $\langle proof \rangle$

**derive-generic-setup** *monoidl*  
 $\langle proof \rangle$

**derive-generic-setup** *group*  
 $\langle proof \rangle$

**instantiation** *int and unit:: semigroup*  
**begin**  
**definition** *mult-int-def* : *mult (x::int) y = x + y*  
**definition** *mult-unit-def*: *mult (x::unit) y = x*  
**instance**  $\langle proof \rangle$   
**end**  
**instantiation** *int and unit:: monoidl*  
**begin**  
**definition** *neutral-int-def* : *neutral = (0::int)*  
**definition** *neutral-unit-def*: *neutral = ()*  
**instance**  $\langle proof \rangle$   
**end**

```

instantiation int and unit:: group
begin
  definition inverse-int-def : inverse (i::int) = 1 - i
  definition inverse-unit-def: inverse u = ()
instance <proof>
end

instantiation prod and sum :: (semigroup, semigroup) semigroup
begin
  definition mult-prod-def:  $x \otimes y = (\text{fst } x \otimes \text{fst } y, \text{snd } x \otimes \text{snd } y)$ 
  definition mult-sum-def:  $x \otimes y = (\text{case } x \text{ of } \text{Inl } a \Rightarrow (\text{case } y \text{ of } \text{Inl } b \Rightarrow \text{Inl } (a \otimes b) \mid \text{Inr } b \Rightarrow \text{Inr } b) \mid \text{Inr } a \Rightarrow (\text{case } y \text{ of } \text{Inl } b \Rightarrow \text{Inr } a \mid \text{Inr } b \Rightarrow \text{Inr } (a \otimes b)))$ 
instance <proof>
end

instantiation prod and sum :: (monoidl, monoidl) monoidl
begin
  definition neutral-prod-def: neutral = (neutral,neutral)
  definition neutral-sum-def: neutral = Inl neutral
instance <proof>
end

instantiation prod :: (group, group) group
begin
  definition inverse-prod-def: inverse p = (inverse (fst p), inverse (snd p))
instance <proof>
end

derive-generic semigroup simple-int <proof>
derive-generic monoidl simple-int <proof>

derive-generic semigroup either <proof>
derive-generic monoidl either <proof>

lemma (B 1 6)  $\otimes$  (B 4 5) = B 4 11 <proof>
lemma (A 2)  $\otimes$  (A 3) = A 5 <proof>
lemma (B 1 6)  $\otimes$  1 = B 0 6 <proof>

lemma (L 3)  $\otimes$  ((L 4)::(int,int) either) = L 7 <proof>
lemma (R (2::int))  $\otimes$  (L (3::int)) = R 2 <proof>

```



**derive-generic** *semigroup list*  
⟨*proof*⟩

**derive-generic** *semigroup tree*  
⟨*proof*⟩

**derive-generic** *monoidl list*  
⟨*proof*⟩

**derive-generic** *monoidl tree*  
⟨*proof*⟩

**lemma**  $[1,2,3,4::int] \otimes [1,2,3] = [2,4,6,4]$  ⟨*proof*⟩

**lemma**  $(Node (3::int) Leaf Leaf) \otimes (Node (1::int) Leaf Leaf) = (Node 4 Leaf Leaf)$  ⟨*proof*⟩

**end**

## References

- [1] J. P. Magalhães, A. Dijkstra, J. Jeuring, and A. Löh. A generic deriving mechanism for Haskell. *ACM Sigplan Notices*, 45(11):37–48, 2010.
- [2] Miles Sabin. shapeless: generic programming for Scala. <https://github.com/milessabin/shapeless>, 2018. [Online; accessed 17-April-2018].