

Descartes' Rule of Signs

Manuel Eberl

December 7, 2022

Abstract

In this work, we formally proved Descartes Rule of Signs, which relates the number of positive real roots of a polynomial with the number of sign changes in its coefficient list.

Our proof follows the simple inductive proof given by Arthan [1], which was also used by John Harrison in his HOL Light formalisation. We proved most of the lemmas for arbitrary linearly-ordered integrity domains (e.g. integers, rationals, reals); the main result, however, requires the intermediate value theorem and was therefore only proven for real polynomials.

Contents

1	Sign changes and Descartes' Rule of Signs	1
1.1	Polynomials	2
1.2	List of partial sums	3
1.3	Sign changes in a list	4
1.4	Arthan's lemma	7
1.5	Roots of a polynomial with a certain property	11
1.6	Coefficient sign changes of a polynomial	13
1.7	Proof of Descartes' sign rule	15

1 Sign changes and Descartes' Rule of Signs

theory *Descartes-Sign-Rule*

imports

Complex-Main

HOL-Computational-Algebra.Polynomial

begin

lemma *op-plus-0*: $((+) (0 :: 'a :: monoid-add)) = id$

by *auto*

lemma *filter-dropWhile*:

$filter (\lambda x. \neg P x) (dropWhile P xs) = filter (\lambda x. \neg P x) xs$

by (*induction xs simp-all*)

1.1 Polynomials

A real polynomial whose leading and constant coefficients have opposite non-zero signs must have a positive root.

lemma *pos-root-exI*:

assumes $\text{poly } p \ 0 * \text{lead-coeff } p < (0 :: \text{real})$

obtains x **where** $x > 0$ $\text{poly } p \ x = 0$

proof –

have $P: \exists x > 0. \text{poly } p \ x = (0 :: \text{real})$ **if** $\text{lead-coeff } p > 0$ $\text{poly } p \ 0 < 0$ **for** p

proof –

note $\text{that}(1)$

also from $\text{poly-pinfty-gt-lc}[OF \ \langle \text{lead-coeff } p > 0 \rangle]$ **obtain** $x0$

where $\bigwedge x. x \geq x0 \implies \text{poly } p \ x \geq \text{lead-coeff } p$ **by** *auto*

hence $\text{poly } p \ (\text{max } x0 \ 1) \geq \text{lead-coeff } p$ **by** *auto*

finally have $\text{poly } p \ (\text{max } x0 \ 1) > 0$.

with that have $\exists x. x > 0 \wedge x < \text{max } x0 \ 1 \wedge \text{poly } p \ x = 0$

by (*intro poly-IVT mult-neg-pos*) *auto*

thus $\exists x > 0. \text{poly } p \ x = 0$ **by** *auto*

qed

show *?thesis*

proof (*cases lead-coeff p > 0*)

case *True*

with *assms* **have** $\text{poly } p \ 0 < 0$

by (*auto simp: mult-less-0-iff*)

from $P[OF \ \text{True} \ \text{this}]$ **that** **show** *?thesis*

by *blast*

next

case *False*

from *False assms* **have** $\text{poly } (-p) \ 0 < 0$

by (*auto simp: mult-less-0-iff*)

moreover from *assms* **have** $p \neq 0$

by *auto*

with *False* **have** $\text{lead-coeff } (-p) > 0$

by (*cases rule: linorder-cases[of lead-coeff p 0]*)

(*simp-all add:*)

ultimately show *?thesis* **using that** $P[\text{of } -p]$ **by** *auto*

qed

qed

Substitute X with aX in a polynomial $p(X)$. This turns all the $X - a$ factors in p into factors of the form $X - 1$.

definition *reduce-root* **where**

$\text{reduce-root } a \ p = \text{pcompose } p \ [:\ 0, \ a:]$

lemma *reduce-root-pCons*:

$\text{reduce-root } a \ (\text{pCons } c \ p) = \text{pCons } c \ (\text{smult } a \ (\text{reduce-root } a \ p))$

by (*simp add: reduce-root-def pcompose-pCons*)

lemma *reduce-root-nonzero* [*simp*]:
 $a \neq 0 \implies p \neq 0 \implies \text{reduce-root } a \ p \neq (0 :: 'a :: \text{idom poly})$
unfolding *reduce-root-def* **using** *pcompose-eq-0*[*of p [:0, a:]*]
by *auto*

1.2 List of partial sums

We first define, for a given list, the list of accumulated partial sums from left to right: the list *psums xs* has as its *i*-th entry $\sum_{j=0}^i xs_j$.

fun *psums* **where**
psums [] = []
| *psums* [x] = [x]
| *psums* (x#y#xs) = x # *psums* ((x+y) # xs)

lemma *length-psums* [*simp*]: $\text{length } (\text{psums } xs) = \text{length } xs$
by (*induction xs rule: psums.induct*) *simp-all*

lemma *psums-Cons*:
 $\text{psums } (x\#xs) = (x :: 'a :: \text{semigroup-add}) \# \text{map } ((+) \ x) \ (\text{psums } xs)$
by (*induction xs rule: psums.induct*) (*simp-all add: algebra-simps*)

lemma *last-psums*:
 $(xs :: 'a :: \text{monoid-add list}) \neq [] \implies \text{last } (\text{psums } xs) = \text{sum-list } xs$
by (*induction xs rule: psums.induct*)
(*auto simp add: add.assoc [symmetric] psums-Cons o-def*)

lemma *psums-0-Cons* [*simp*]:
 $\text{psums } (0\#xs :: 'a :: \text{monoid-add list}) = 0 \# \text{psums } xs$
by (*induction xs rule: psums.induct*) (*simp-all add: algebra-simps*)

lemma *map-uminus-psums*:
fixes $xs :: 'a :: \text{ab-group-add list}$
shows $\text{map } \text{uminus } (\text{psums } xs) = \text{psums } (\text{map } \text{uminus } xs)$
by (*induction xs rule: psums.induct*) (*simp-all*)

lemma *psums-replicate-0-append*:
 $\text{psums } (\text{replicate } n \ (0 :: 'a :: \text{monoid-add}) \ @ \ xs) =$
 $\text{replicate } n \ 0 \ @ \ \text{psums } xs$
by (*induction n*) (*simp-all add: psums-Cons op-plus-0*)

lemma *psums-nth*: $n < \text{length } xs \implies \text{psums } xs \ ! \ n = (\sum_{i \leq n} xs \ ! \ i)$
proof (*induction xs arbitrary: n rule: psums.induct[case-names Nil sng rec]*)
case (*rec x y xs n*)
show ?*case*
proof (*cases n*)
case (*Suc m*)
from *Suc* **have** $\text{psums } (x \ # \ y \ # \ xs) \ ! \ n = \text{psums } ((x+y) \ # \ xs) \ ! \ m$ **by** *simp*
also from *rec.prem*s *Suc* **have** $\dots = (\sum_{i \leq m} ((x+y) \ # \ xs) \ ! \ i)$
by (*intro rec.IH*) *simp-all*

also have $\dots = x + y + (\sum_{i=1..m}. (y\#xs) ! i)$
by (*auto simp: atLeast0AtMost [symmetric] sum.atLeast-Suc-atMost[of 0]*)
also have $(\sum_{i=1..m}. (y\#xs) ! i) = (\sum_{i=Suc\ 1..Suc\ m}. (x\#y\#xs) ! i)$
by (*subst sum.shift-bounds-cl-Suc-ivl simp*)
also from *Suc* **have** $x + y + \dots = (\sum_{i \leq n}. (x\#y\#xs) ! i)$
by (*auto simp: atLeast0AtMost [symmetric] sum.atLeast-Suc-atMost add-ac*)
finally show *?thesis* .
qed *simp*
qed *simp-all*

1.3 Sign changes in a list

Next, we define the number of sign changes in a sequence. Intuitively, this is the number of times that, when passing through the list, a sign change between one element and the next element occurs (while ignoring all zero entries).

We implement this by filtering all zeros from the list of signs, removing all adjacent equal elements and taking the length of the resulting list minus one.

definition *sign-changes* :: (*'a* :: {*sgn,zero*} *list*) \Rightarrow *nat* **where**
*sign-changes xs = length (remdups-adj (filter ($\lambda x. x \neq 0$) (map *sgn xs*))) - 1*

lemma *sign-changes-Nil* [*simp*]: *sign-changes [] = 0*
by (*simp add: sign-changes-def*)

lemma *sign-changes-singleton* [*simp*]: *sign-changes [x] = 0*
by (*simp add: sign-changes-def*)

lemma *sign-changes-cong*:
assumes *map sgn xs = map sgn ys*
shows *sign-changes xs = sign-changes ys*
using *assms unfolding sign-changes-def by simp*

lemma *sign-changes-Cons-ge*: *sign-changes (x # xs) \geq sign-changes xs*
unfolding *sign-changes-def by (simp add: remdups-adj-Cons split: list.split)*

lemma *sign-changes-Cons-Cons-different*:
fixes *x y :: 'a :: linordered-idom*
assumes *x * y < 0*
shows *sign-changes (x # y # xs) = 1 + sign-changes (y # xs)*

proof –
from *assms* **have** *sgn x = -1 \wedge sgn y = 1 \vee sgn x = 1 \wedge sgn y = -1*
by (*auto simp: mult-less-0-iff*)
thus *?thesis* **by** (*fastforce simp: sign-changes-def*)
qed

lemma *sign-changes-Cons-Cons-same*:
fixes *x y :: 'a :: linordered-idom*

shows $x * y > 0 \implies \text{sign-changes } (x \# y \# xs) = \text{sign-changes } (y \# xs)$
by (*subst (asm) zero-less-mult-iff*) (*fastforce simp: sign-changes-def*)

lemma *sign-changes-0-Cons* [*simp*]:
 $\text{sign-changes } (0 \# xs :: 'a :: \text{idom-abs-sgn list}) = \text{sign-changes } xs$
by (*simp add: sign-changes-def*)

lemma *sign-changes-two*:
fixes $x y :: 'a :: \text{linordered-idom}$
shows $\text{sign-changes } [x,y] =$
 $(\text{if } x > 0 \wedge y < 0 \vee x < 0 \wedge y > 0 \text{ then } 1 \text{ else } 0)$
by (*auto simp: sgn-if sign-changes-def mult-less-0-iff*)

lemma *sign-changes-induct* [*case-names nil sing zero nonzero*]:
assumes $P [] \wedge x. P [x] \wedge xs. P xs \implies P (0 \# xs)$
 $\wedge x y xs. x \neq 0 \implies P ((x + y) \# xs) \implies P (x \# y \# xs)$
shows $P xs$
proof (*induction length xs arbitrary: xs rule: less-induct*)
case (*less xs*)
show *?case*
proof (*cases xs rule: psums.cases*)
fix $x y xs'$ **assume** $xs = x \# y \# xs'$
with *assms less* **show** *?thesis* **by** (*cases x = 0*) *auto*
qed (*insert less assms, auto*)

qed

lemma *sign-changes-filter*:
fixes $xs :: 'a :: \text{linordered-idom list}$
shows $\text{sign-changes } (\text{filter } (\lambda x. x \neq 0) xs) = \text{sign-changes } xs$
by (*simp add: sign-changes-def filter-map o-def sgn-0-0*)

lemma *sign-changes-Cons-Cons-0*:
fixes $xs :: 'a :: \text{linordered-idom list}$
shows $\text{sign-changes } (x \# 0 \# xs) = \text{sign-changes } (x \# xs)$
by (*subst (1 2) sign-changes-filter [symmetric] simp-all*)

lemma *sign-changes-uminus*:
fixes $xs :: 'a :: \text{linordered-idom list}$
shows $\text{sign-changes } (\text{map } \text{uminus } xs) = \text{sign-changes } xs$
proof –
have $\text{sign-changes } (\text{map } \text{uminus } xs) =$
 $\text{length } (\text{remdups-adj } [x \leftarrow \text{map } \text{sgn } (\text{map } \text{uminus } xs) . x \neq 0]) - 1$
unfolding *sign-changes-def ..*
also have $\text{map } \text{sgn } (\text{map } \text{uminus } xs) = \text{map } \text{uminus } (\text{map } \text{sgn } xs)$
by (*auto simp: sgn-minus*)
also have $\text{remdups-adj } (\text{filter } (\lambda x. x \neq 0) \dots) =$
 $\text{map } \text{uminus } (\text{remdups-adj } (\text{filter } (\lambda x. x \neq 0) (\text{map } \text{sgn } xs)))$
by (*subst filter-map, subst remdups-adj-map-injective*)
(simp-all add: o-def)

also have $\text{length } \dots - 1 = \text{sign-changes } xs$ by (simp add: sign-changes-def)
 finally show ?thesis .
 qed

lemma *sign-changes-replicate*: $\text{sign-changes } (\text{replicate } n \ x) = 0$
 by (simp add: sign-changes-def remdups-adj-replicate filter-replicate)

lemma *sign-changes-decompose*:
 assumes $x \neq (0 :: 'a :: \text{linordered-idom})$
 shows $\text{sign-changes } (xs \ @ \ x \ \# \ ys) =$
 $\text{sign-changes } (xs \ @ \ [x]) + \text{sign-changes } (x \ \# \ ys)$

proof –
 have $\text{sign-changes } (xs \ @ \ x \ \# \ ys) =$
 $\text{length } (\text{remdups-adj } ([x \leftarrow \text{map } \text{sgn } xs . x \neq 0] \ @$
 $\text{sgn } x \ \# \ [x \leftarrow \text{map } \text{sgn } ys . x \neq 0])) - 1$
 by (simp add: sgn-0-0 assms sign-changes-def)
 also have $\dots = \text{sign-changes } (xs \ @ \ [x]) + \text{sign-changes } (x \ \# \ ys)$
 by (subst remdups-adj-append) (simp add: sign-changes-def assms sgn-0-0)
 finally show ?thesis .
 qed

If the first and the last entry of a list are non-zero, its number of sign changes is even if and only if the first and the last element have the same sign. This will be important later to establish the base case of Descartes' Rule. (if there are no positive roots, the number of sign changes is even)

lemma *even-sign-changes-iff*:
 assumes $xs \neq ([] :: 'a :: \text{linordered-idom list})$ $\text{hd } xs \neq 0$ $\text{last } xs \neq 0$
 shows $\text{even } (\text{sign-changes } xs) \longleftrightarrow \text{sgn } (\text{hd } xs) = \text{sgn } (\text{last } xs)$
 using assms
proof (induction length xs arbitrary: xs rule: less-induct)
 case (less xs)
 show ?case
proof (cases xs)
 case (Cons x xs')
 note $x = \text{this}$
 show ?thesis
proof (cases xs')
 case (Cons y xs'')
 note $y = \text{this}$
 show ?thesis
proof (rule linorder-cases[of $x*y$ 0])
 assume $xy: x*y = 0$
 with $x \ y$ less(1,3,4) show ?thesis by (auto simp: sign-changes-Cons-Cons-0)
 next
 assume $xy: x*y > 0$
 with less(1,4) show ?thesis
 by (auto simp add: $x \ y$ sign-changes-Cons-Cons-same zero-less-mult-iff)
 next
 assume $xy: x*y < 0$

```

moreover from  $xy$  have  $sgn\ x = -\ sgn\ y$  by (auto simp: mult-less-0-iff)
moreover have even (sign-changes ( $y \# xs''$ ))  $\longleftrightarrow$ 
       $sgn\ (hd\ (y \# xs'')) = sgn\ (last\ (y \# xs''))$ 
  using  $xy$  less.prems by (intro less) (auto simp: x y)
moreover from  $xy$  less.prems
  have  $sgn\ y = sgn\ (last\ xs) \longleftrightarrow -sgn\ y \neq sgn\ (last\ xs)$ 
  by (auto simp: sgn-if)
ultimately show ?thesis by (auto simp: sign-changes-Cons-Cons-different
 $x\ y$ )
  qed
qed (auto simp: x)
qed (insert less.prems, simp-all)
qed

```

1.4 Arthan's lemma

```

context
begin

```

We first prove an auxiliary lemma that allows us to assume w.l.o.g. that the first element of the list is non-negative, similarly to what Arthan does in his proof.

```

private lemma arthan-wlog [consumes 3, case-names nonneg lift]:
  fixes  $xs :: 'a :: linordered-idom\ list$ 
  assumes  $xs \neq []$   $last\ xs \neq 0$   $x + y + sum-list\ xs = 0$ 
  assumes  $\bigwedge x\ y\ xs. xs \neq [] \implies last\ xs \neq 0 \implies$ 
       $x + y + sum-list\ xs = 0 \implies x \geq 0 \implies P\ x\ y\ xs$ 
  assumes  $\bigwedge x\ y\ xs. xs \neq [] \implies P\ x\ y\ xs \implies P\ (-x)\ (-y)\ (map\ uminus\ xs)$ 
  shows  $P\ x\ y\ xs$ 
proof (cases  $x \geq 0$ )
  assume  $x: \neg(x \geq 0)$ 
  from assms have  $map\ uminus\ xs \neq []$  by simp
  moreover from  $x$  assms(1,2,3) have  $P\ (-x)\ (-y)\ (map\ uminus\ xs)$ 
  using uminus-sum-list-map[of  $\lambda x. x\ xs$ , symmetric]
  by (intro assms) (auto simp: last-map algebra-simps o-def neg-eq-iff-add-eq-0)
  ultimately have  $P\ (-(-x))\ (-(-y))\ (map\ uminus\ (map\ uminus\ xs))$  by (rule
assms)
  thus ?thesis by (simp add: o-def)
qed (simp-all add: assms)

```

We now show that the α and β in Arthan's proof have the necessary properties: their difference is non-negative and even.

```

private lemma arthan-aux1:
  fixes  $xs :: 'a :: \{linordered-idom\}\ list$ 
  assumes  $xs \neq []$   $last\ xs \neq 0$   $x + y + sum-list\ xs = 0$ 
  defines  $v \equiv \lambda xs. int\ (sign-changes\ xs)$ 
  shows  $v\ (x \# y \# xs) - v\ ((x + y) \# xs) \geq$ 
       $v\ (psums\ (x \# y \# xs)) - v\ (psums\ ((x + y) \# xs)) \wedge$ 
      even ( $v\ (x \# y \# xs) - v\ ((x + y) \# xs) -$ 

```

```

      (v (psums (x # y # xs)) - v (psums ((x + y) # xs)))
using assms(1-3)
proof (induction rule: arthan-wlog)
  have uminus-v: v (map uminus xs) = v xs for xs by (simp add: v-def sign-changes-uminus)

  case (lift x y xs)
  note lift(2)
  also have v (psums (x#y#xs)) - v (psums ((x+y)#xs)) =
    v (psums (- x # - y # map uminus xs)) -
    v (psums ((- x + - y) # map uminus xs))
    by (subst (1 2) uminus-v [symmetric]) (simp add: map-uminus-psums)
  also have v (x # y # xs) - v ((x + y) # xs) =
    v (-x # -y # map uminus xs) - v ((-x + -y) # map uminus xs)
    by (subst (1 2) uminus-v [symmetric]) simp
  finally show ?case .
next
  case (nonneg x y xs)
  define p where p = (LEAST n. xs ! n ≠ 0)
  define xs1 :: 'a list where xs1 = replicate p 0
  define xs2 where xs2 = drop (Suc p) xs
  from nonneg have xs ! (length xs - 1) ≠ 0 by (simp add: last-conv-nth)
  hence p-nz: xs ! p ≠ 0 unfolding p-def by (rule LeastI)
  {
    fix q assume q < p hence xs ! q = 0
    using Least-le[of λn. xs ! n ≠ 0 q] unfolding p-def by force
  } note less-p-zero = this
  from Least-le[of λn. xs ! n ≠ 0 length xs - 1] nonneg
    have p ≤ length xs - 1 unfolding p-def by (auto simp: last-conv-nth)
  with nonneg have p-less-length: p < length xs by (cases xs) simp-all

  from p-less-length less-p-zero have take p xs = replicate p 0
    by (subst list-eq-iff-nth-eq) auto
  with p-less-length have xs-decompose: xs = xs1 @ xs ! p # xs2
    unfolding xs1-def xs2-def
    by (subst append-take-drop-id [of p, symmetric],
      subst Cons-nth-drop-Suc) simp-all

  have v-decompose: v (xs' @ xs) = v (xs' @ [xs ! p]) + v (xs ! p # xs2) for xs'
  proof -
    have xs' @ xs = (xs' @ xs1) @ xs ! p # xs2 by (subst xs-decompose) simp
    also have v ... = v (xs' @ [xs ! p]) + v (xs ! p # xs2) unfolding v-def
      by (subst sign-changes-decompose[OF p-nz],
        subst (1 2 3 4) sign-changes-filter [symmetric]) (simp-all add: xs1-def)
    finally show ?thesis .
  qed

  have psums-decompose: psums xs = replicate p 0 @ psums (xs!p # xs2)
    by (subst xs-decompose) (simp add: xs1-def psums-replicate-0-append)
  have v-psums-decompose: sign-changes (xs' @ psums xs) = sign-changes (xs' @

```



```

[xs!p] +
  sign-changes (xs!p # map ((+) (xs!p)) (psums xs2)) for xs'
proof -
  fix xs' :: 'a list
  have sign-changes (xs' @ psums xs) =
    sign-changes (xs' @ xs ! p # map ((+) (xs!p)) (psums xs2))
  by (subst psums-decompose, subst (1 2) sign-changes-filter [symmetric])
    (simp-all add: psums-Cons)
  also have ... = sign-changes (xs' @ [xs!p]) +
    sign-changes (xs!p # map ((+) (xs!p)) (psums xs2))
  by (subst sign-changes-decompose[OF p-nz]) simp-all
  finally show sign-changes (xs' @ psums xs) = ... .
qed

show ?case
proof (cases x > 0)
  assume ¬(x > 0)
  with nonneg show ?thesis by (auto simp: v-def)
next
  assume x: x > 0
  show ?thesis
  proof (rule linorder-cases[of y 0])
    assume y: y > 0
    from x and this have xy: x + y > 0 by (rule add-pos-pos)
    with y have sign-changes ((x + y) # xs) = sign-changes (y # xs)
    by (intro sign-changes-cong) auto
    moreover have sign-changes (x # psums ((x + y) # xs)) =
      sign-changes (psums ((x+y) # xs))
    using xy by (subst (1 2) psums-Cons) (simp-all add: sign-changes-Cons-Cons-same)
    ultimately show ?thesis using x y
    by (simp add: v-def algebra-simps sign-changes-Cons-Cons-same)
  next
    assume y: y = 0
    with x show ?thesis
    by (simp add: v-def sign-changes-Cons-Cons-0 psums-Cons
      o-def sign-changes-Cons-Cons-same)
  next
    assume y: y < 0
    with x have different: x * y < 0 by (rule mult-pos-neg)
    show ?thesis
    proof (rule linorder-cases[of x + y 0])
      assume xy: x + y < 0
      with x have different': x * (x + y) < 0 by (rule mult-pos-neg)
      have ( $\lambda t. t + (x + y) = ((+) (x + y))$ ) by (rule ext) simp
      moreover from y xy have sign-changes ((x+y) # xs) = sign-changes (y
# xs)
      by (intro sign-changes-cong) auto
      ultimately show ?thesis using xy different different' y
      by (simp add: v-def sign-changes-Cons-Cons-different psums-Cons o-def

```

```

add-ac)
next
  assume  $xy: x + y = 0$ 
  show ?case
  proof (cases  $xs ! p > 0$ )
    assume  $p: xs ! p > 0$ 
    from  $p y$  have  $different'$ :  $y * xs ! p < 0$  by (intro mult-neg-pos)
    with  $v-decompose[of [x, y]] v-decompose[of [x+y]] x xy p different different'$ 

       $v-psums-decompose[of [x]] v-psums-decompose[of []]$ 
    show ?thesis by (auto simp add: algebra-simps v-def sign-changes-Cons-Cons-0

      sign-changes-Cons-Cons-different sign-changes-Cons-Cons-same)

  next
    assume  $\neg(xs ! p > 0)$ 
    with  $p-nz$  have  $p: xs ! p < 0$  by simp
    from  $p y$  have  $same: y * xs ! p > 0$  by (intro mult-neg-neg)
    from  $p x$  have  $different'$ :  $x * xs ! p < 0$  by (intro mult-pos-neg)
    from  $v-decompose[of [x, y]] v-decompose[of [x+y]] xy different different'$ 
  same
     $v-psums-decompose[of [x]] v-psums-decompose[of []]$ 
  show ?thesis by (auto simp add: algebra-simps v-def sign-changes-Cons-Cons-0

    sign-changes-Cons-Cons-different sign-changes-Cons-Cons-same)

qed
next
  assume  $xy: x + y > 0$ 
  from  $x$  and  $this$  have  $same: x * (x + y) > 0$  by (rule mult-pos-pos)
  show ?case
  proof (cases  $xs ! p > 0$ )
    assume  $p: xs ! p > 0$ 
    from  $xy p$  have  $same'$ :  $(x + y) * xs ! p > 0$  by (intro mult-pos-pos)
    from  $p y$  have  $different'$ :  $y * xs ! p < 0$  by (intro mult-neg-pos)
    have  $(\lambda t. t + (x + y)) = ((+) (x + y))$  by (rule ext) simp
    with  $v-decompose[of [x, y]] v-decompose[of [x+y]] different different' same$ 
  same'
    show ?thesis by (auto simp add: algebra-simps v-def psums-Cons o-def

      sign-changes-Cons-Cons-different sign-changes-Cons-Cons-same)

  next
    assume  $\neg(xs ! p > 0)$ 
    with  $p-nz$  have  $p: xs ! p < 0$  by simp
    from  $xy p$  have  $different'$ :  $(x + y) * xs ! p < 0$  by (rule mult-pos-neg)
    from  $y p$  have  $same'$ :  $y * xs ! p > 0$  by (rule mult-neg-neg)
    have  $(\lambda t. t + (x + y)) = ((+) (x + y))$  by (rule ext) simp
    with  $v-decompose[of [x, y]] v-decompose[of [x+y]] different different' same$ 
  same'
    show ?thesis by (auto simp add: algebra-simps v-def psums-Cons o-def

      sign-changes-Cons-Cons-different sign-changes-Cons-Cons-same)

qed

```

qed
 qed
 qed
 qed

Now we can prove the main lemma of the proof by induction over the list with our specialised induction rule for *sign-changes*. It states that for a non-empty list whose last element is non-zero and whose sum is zero, the difference of the sign changes in the list and in the list of its partial sums is odd and positive.

lemma *arthan*:
fixes $xs :: 'a :: \text{linordered-idom list}$
assumes $xs \neq []$ $\text{last } xs \neq 0$ $\text{sum-list } xs = 0$
shows $\text{sign-changes } xs > \text{sign-changes } (\text{psums } xs) \wedge$
 $\text{odd } (\text{sign-changes } xs - \text{sign-changes } (\text{psums } xs))$
using *assms*
proof (*induction xs rule: sign-changes-induct*)
case (*nonzero x y xs*)
show ?*case*
proof (*cases xs = []*)
case *False*
define α **where** $\alpha = \text{int } (\text{sign-changes } (x \# y \# xs)) - \text{int } (\text{sign-changes } ((x + y) \# xs))$
define β **where** $\beta = \text{int } (\text{sign-changes } (\text{psums } (x \# y \# xs))) - \text{int } (\text{sign-changes } (\text{psums } ((x+y) \# xs)))$
from *nonzero False* **have** $\alpha \geq \beta \wedge \text{even } (\alpha - \beta)$ **unfolding** $\alpha\text{-def } \beta\text{-def}$
by (*intro arthan-aux1*) *auto*
from *False* **and** *nonzero.prem*s **have**
 $\text{sign-changes } (\text{psums } ((x + y) \# xs)) < \text{sign-changes } ((x + y) \# xs) \wedge$
 $\text{odd } (\text{sign-changes } ((x + y) \# xs) - \text{sign-changes } (\text{psums } ((x + y) \# xs)))$
by (*intro nonzero.IH*) (*auto simp: add.assoc*)
with *arthan-aux1* [*of xs x y*] *nonzero(4,5)* *False(1)* **show** ?*thesis* **by** *force*
qed (*insert nonzero.prem*s, *auto split: if-split-asm simp: sign-changes-two add-eq-0-iff*)
qed (*auto split: if-split-asm simp: add-eq-0-iff*)
end

1.5 Roots of a polynomial with a certain property

The set of roots of a polynomial p that fulfil a given property P :

definition *roots-with* $P p = \{x. P x \wedge \text{poly } p x = 0\}$

The number of roots of a polynomial p with a given property P , where multiple roots are counted multiple times.

definition *count-roots-with* $P p = (\sum_{x \in \text{roots-with } P p. \text{order } x p}$

abbreviation *pos-roots* $\equiv \text{roots-with } (\lambda x. x > 0)$

abbreviation *count-pos-roots* $\equiv \text{count-roots-with } (\lambda x. x > 0)$

lemma *finite-roots-with* [simp]:
 ($p :: 'a :: \text{linordered-idom poly} \neq 0 \implies \text{finite} (\text{roots-with } P p)$)
by (rule *finite-subset*[OF - *poly-roots-finite*[of p]]) (auto simp: *roots-with-def*)

lemma *count-roots-with-times-root*:
assumes $p \neq 0$ $P (a :: 'a :: \text{linordered-idom})$
shows $\text{count-roots-with } P ([:a, -1:] * p) = \text{Suc} (\text{count-roots-with } P p)$
proof –
define q **where** $q = [:a, -1:] * p$
from *assms* **have** $a \in \text{roots-with } P q$ **by** (simp-all add: *roots-with-def* $q\text{-def}$)
have $q\text{-nz}$: $q \neq 0$ **unfolding** $q\text{-def}$ **by** (rule *no-zero-divisors*) (simp-all add: *assms*)

have $\text{count-roots-with } P q = (\sum_{x \in \text{roots-with } P q} \text{order } x q)$ **by** (simp add: *count-roots-with-def*)
also from $q\text{-nz}$ **have** $\dots = \text{order } a q + (\sum_{x \in \text{roots-with } P q - \{a\}} \text{order } x q)$
by (*subst sum.remove*) *simp-all*
also have $\text{order } a q = \text{order } a [:a, -1:] + \text{order } a p$ **unfolding** $q\text{-def}$
by (*subst order-mult*[OF *no-zero-divisors*]) (simp-all add: *assms*)
also have $\text{order } a [:a, -1:] = 1$
by (*subst order-smult* [of -1 , *symmetric*])
 (*insert order-power-n-n*[of a 1], *simp-all* add: *order-1*)
also have $(\sum_{x \in \text{roots-with } P q - \{a\}} \text{order } x q) = (\sum_{x \in \text{roots-with } P q - \{a\}} \text{order } x p)$
order x p
proof (*intro sum.cong refl*)
fix x **assume** $x \in \text{roots-with } P q - \{a\}$
from *assms* **have** $\text{order } x q = \text{order } x [:a, -1:] + \text{order } x p$ **unfolding** $q\text{-def}$
by (*subst order-mult*[OF *no-zero-divisors*]) (simp-all add: *assms*)
also from x **have** $\text{order } x [:a, -1:] = 0$ **by** (*intro order-0I*) *simp-all*
finally show $\text{order } x q = \text{order } x p$ **by** *simp*
qed
also from $q\text{-nz}$ **have** $1 + \text{order } a p + (\sum_{x \in \text{roots-with } P q - \{a\}} \text{order } x p)$
 =
 $1 + (\sum_{x \in \text{roots-with } P q} \text{order } x p)$
by (*subst add.assoc*, *subst sum.remove*[*symmetric*]) *simp-all*
also from $q\text{-nz}$ **have** $(\sum_{x \in \text{roots-with } P q} \text{order } x p) = (\sum_{x \in \text{roots-with } P p} \text{order } x p)$
order x p
proof (*intro sum.mono-neutral-right*)
show $\text{roots-with } P p \subseteq \text{roots-with } P q$
by (auto simp: *roots-with-def* $q\text{-def}$ *simp del*: *mult-pCons-left*)
show $\forall x \in \text{roots-with } P q - \text{roots-with } P p. \text{order } x p = 0$
by (auto simp: *roots-with-def* $q\text{-def}$ *order-root* *simp del*: *mult-pCons-left*)
qed *simp-all*
finally show *?thesis* **by** (simp add: $q\text{-def}$ *count-roots-with-def*)
qed

1.6 Coefficient sign changes of a polynomial

abbreviation (*input*) $\text{coeff-sign-changes } f \equiv \text{sign-changes } (\text{coeffs } f)$

We first show that when building a polynomial from a coefficient list, the coefficient sign sign changes of the resulting polynomial are the same as the same sign changes in the list.

Note that constructing a polynomial from a list removes all trailing zeros.

lemma *sign-changes-coeff-sign-changes*:

assumes $\text{Poly } xs = (p :: 'a :: \text{linordered-idom } \text{poly})$

shows $\text{sign-changes } xs = \text{coeff-sign-changes } p$

proof –

have $\text{coeffs } p = \text{coeffs } (\text{Poly } xs)$ **by** (*subst assms*) (*rule refl*)

also have $\dots = \text{strip-while } ((=) 0) xs$ **by** *simp*

also have $\text{filter } ((\neq) 0) \dots = \text{filter } ((\neq) 0) xs$ **unfolding** *strip-while-def o-def*

by (*subst rev-filter [symmetric]*, *subst filter-dropWhile*) (*simp-all add: rev-filter*)

also have $\text{sign-changes } \dots = \text{sign-changes } xs$ **by** (*simp add: sign-changes-filter*)

finally show *?thesis* **by** (*simp add: sign-changes-filter*)

qed

By applying *reduce-root a*, we can assume w.l.o.g. that the root in question is 1, since applying root reduction does not change the number of sign changes.

lemma *coeff-sign-changes-reduce-root*:

assumes $a > (0 :: 'a :: \text{linordered-idom})$

shows $\text{coeff-sign-changes } (\text{reduce-root } a p) = \text{coeff-sign-changes } p$

proof (*intro sign-changes-cong, induction p*)

case (*pCons c p*)

have $\text{map } \text{sgn } (\text{coeffs } (\text{reduce-root } a (pCons c p))) =$

$cCons (\text{sgn } c) (\text{map } \text{sgn } (\text{coeffs } (\text{reduce-root } a p)))$

using *assms* **by** (*auto simp add: cCons-def sgn-0-0 sgn-mult reduce-root-pCons coeffs-smult*)

also note *pCons.IH*

also have $cCons (\text{sgn } c) (\text{map } \text{sgn } (\text{coeffs } p)) = \text{map } \text{sgn } (\text{coeffs } (pCons c p))$

using *assms* **by** (*auto simp add: cCons-def sgn-0-0*)

finally show *?case* .

qed (*simp-all add: reduce-root-def*)

Multiplying a polynomial with a positive constant also does not change the number of sign changes. (in fact, any non-zero constant would also work, but the proof is slightly more difficult and positive constants suffice in our use case)

lemma *coeff-sign-changes-smult*:

assumes $a > (0 :: 'a :: \text{linordered-idom})$

shows $\text{coeff-sign-changes } (\text{smult } a p) = \text{coeff-sign-changes } p$

using *assms* **by** (*auto intro!: sign-changes-cong simp: sgn-mult coeffs-smult*)

context

begin

We now show that a polynomial with an odd number of sign changes contains a positive root. We first assume that the constant coefficient is non-zero. Then it is clear that the polynomial's sign at 0 will be the sign of the constant coefficient, whereas the polynomial's sign for sufficiently large inputs will be the sign of the leading coefficient.

Moreover, we have shown before that in a list with an odd number of sign changes and non-zero initial and last coefficients, the initial coefficient and the last coefficient have opposite and non-zero signs. Then, the polynomial obviously has a positive root.

private lemma *odd-coeff-sign-changes-imp-pos-roots-aux*:

assumes *[simp]*: $p \neq (0 :: \text{real poly})$ *poly* p $0 \neq 0$

assumes *odd* (*coeff-sign-changes* p)

obtains x **where** $x > 0$ *poly* p $x = 0$

proof –

from $\langle \text{poly } p \ 0 \neq 0 \rangle$

have *[simp]*: $\text{hd} (\text{coeffs } p) \neq 0$

by (*induct* p) *auto*

from *assms* **have** $\neg \text{even} (\text{coeff-sign-changes } p)$

by *blast*

also have $\text{even} (\text{coeff-sign-changes } p) \longleftrightarrow \text{sgn} (\text{hd} (\text{coeffs } p)) = \text{sgn} (\text{lead-coeff } p)$

by (*auto simp add: even-sign-changes-iff last-coeffs-eq-coeff-degree*)

finally have $\text{sgn} (\text{hd} (\text{coeffs } p)) * \text{sgn} (\text{lead-coeff } p) < 0$

by (*auto simp: sgn-if split: if-split-asm*)

also from $\langle p \neq 0 \rangle$ **have** $\text{hd} (\text{coeffs } p) = \text{poly } p \ 0$ **by** (*induction* p) *auto*

finally have $\text{poly } p \ 0 * \text{lead-coeff } p < 0$ **by** (*auto simp: mult-less-0-iff*)

from *pos-root-exI*[*OF this*] **that show** *?thesis* **by** *blast*

qed

We can now show the statement without the restriction to a non-zero constant coefficient. We can do this by simply factoring p into the form $p \cdot x^n$, where n is chosen as large as possible. This corresponds to stripping all initial zeros of the coefficient list, which obviously changes neither the existence of positive roots nor the number of coefficient sign changes.

lemma *odd-coeff-sign-changes-imp-pos-roots*:

assumes $p \neq (0 :: \text{real poly})$

assumes *odd* (*coeff-sign-changes* p)

obtains x **where** $x > 0$ *poly* p $x = 0$

proof –

define s **where** $s = \text{sgn} (\text{lead-coeff } p)$

define n **where** $n = \text{order } 0 \ p$

define r **where** $r = p \ \text{div} \ [;0, 1:] \wedge^n$

have $p: p = [;0, 1:] \wedge^n * r$ **unfolding** *r-def n-def*

using *order-1*[*of 0 p*] **by** (*simp del: mult-pCons-left*)

```

from assms p have r-nz:  $r \neq 0$  by auto

obtain x where  $x > 0$  poly  $r\ x = 0$ 
proof (rule odd-coeff-sign-changes-imp-pos-roots-aux)
  show  $r \neq 0$  by fact
  have order 0 p = order 0 p + order 0 r
    by (subst p, insert order-power-n-n[of 0::real n] r-nz)
      (simp del: mult-pCons-left add: order-mult n-def)
  hence order 0 r = 0 by simp
  with r-nz show nz: poly r 0  $\neq 0$  by (simp add: order-root)

  note  $\langle \text{odd } (\text{coeff-sign-changes } p) \rangle$ 
  also have  $p = [:0, 1:] \wedge^n * r$  by (simp add: p)
  also have  $[:0, 1:] \wedge^n = \text{monom } 1\ n$ 
    by (induction n) (simp-all add: monom-Suc monom-0)
  also have coeffs (monom 1 n * r) = replicate n 0 @ coeffs r
    by (induction n) (simp-all add: monom-Suc cCons-def r-nz monom-0)
  also have sign-changes ... = coeff-sign-changes r
    by (subst (1 2) sign-changes-filter [symmetric] simp)
  finally show odd (coeff-sign-changes r) .
qed
thus ?thesis by (intro that[of x] (simp-all add: p))
qed

end

```

1.7 Proof of Descartes' sign rule

For a polynomial $p(X) = a_0 + \dots + a_n X^n$, we have $[X^i](1 - X)p(X) = (\sum_{j=0}^i a_j)$.

```

lemma coeff-poly-times-one-minus-x:
  fixes  $g :: 'a :: \text{linordered-idom } \text{poly}$ 
  shows  $\text{coeff } g\ n = (\sum_{i \leq n}. \text{coeff } (g * [:1, -1:])\ i)$ 
  by (induction n) simp-all

```

We apply the previous lemma to the coefficient list of a polynomial and show: given a polynomial $p(X)$ and $q(X) = (1 - X)p(X)$, the coefficient list of $p(X)$ is the list of partial sums of the coefficient list of $q(X)$.

```

lemma Poly-times-one-minus-x-eq-psums:
  fixes  $xs :: 'a :: \text{linordered-idom } \text{list}$ 
  assumes [simp]:  $\text{length } xs = \text{length } ys$ 
  assumes  $\text{Poly } xs = \text{Poly } ys * [:1, -1:]$ 
  shows  $ys = \text{psums } xs$ 
proof (rule nth-equalityI; safe?)
  fix  $i$  assume  $i < \text{length } ys$ 
  hence  $ys ! i = \text{coeff } (\text{Poly } ys)\ i$ 
  by (simp add: nth-default-def)

```

also from *coeff-poly-times-one-minus-x* [of *Poly ys i*] *assms*
have $\dots = (\sum_{j \leq i} \text{coeff } (\text{Poly } xs) j)$ **by** *simp*
also from *i* **have** $\dots = \text{psums } xs ! i$
by (*auto simp: nth-default-def psums-nth*)
finally show $ys ! i = \text{psums } xs ! i$.
qed *simp-all*

We can now apply our main lemma on the sign changes in lists to the coefficient lists of a nonzero polynomial $p(X)$ and $(1-X)p(X)$: the difference of the changes in the coefficient lists is odd and positive.

lemma *sign-changes-poly-times-one-minus-x*:
fixes $g :: 'a :: \text{linordered-idom poly}$ **and** $a :: 'a$
assumes $nz: g \neq 0$
defines $v \equiv \text{coeff-sign-changes}$
shows $v ([:1, -1:] * g) - v g > 0 \wedge \text{odd } (v ([:1, -1:] * g) - v g)$
proof –
define xs **where** $xs = \text{coeffs } ([:1, -1:] * g)$
define ys **where** $ys = \text{coeffs } g @ [0]$
have $ys: ys = \text{psums } xs$
proof (*rule Poly-times-one-minus-x-eq-psums*)
show $\text{length } xs = \text{length } ys$ **unfolding** *xs-def ys-def*
by (*simp add: length-coeffs nz degree-mult-eq no-zero-divisors del: mult-pCons-left*)
show $\text{Poly } xs = \text{Poly } ys * [:1, -1:]$ **unfolding** *xs-def ys-def*
by (*simp only: Poly-snoc Poly-coeffs*) *simp*
qed
have $\text{sign-changes } (\text{psums } xs) < \text{sign-changes } xs \wedge$
 $\text{odd } (\text{sign-changes } xs - \text{sign-changes } (\text{psums } xs))$
proof (*rule arthan*)
show $xs \neq []$
by (*auto simp: xs-def nz simp del: mult-pCons-left*)
then show $\text{sum-list } xs = 0$ **by** (*simp add: last-psums [symmetric] ys [symmetric]*
ys-def)
show $\text{last } xs \neq 0$
by (*auto simp: xs-def nz last-coeffs-eq-coeff-degree simp del: mult-pCons-left*)
qed
with ys **have** $\text{sign-changes } ys < \text{sign-changes } xs \wedge$
 $\text{odd } (\text{sign-changes } xs - \text{sign-changes } ys)$ **by** *simp*
also have $\text{sign-changes } xs = v ([:1, -1:] * g)$ **unfolding** *v-def*
by (*intro sign-changes-coeff-sign-changes*) (*simp-all add: xs-def*)
also have $\text{sign-changes } ys = v g$ **unfolding** *v-def*
by (*intro sign-changes-coeff-sign-changes*) (*simp-all add: ys-def Poly-snoc*)
finally show *?thesis* **by** *simp*
qed

We can now lift the previous lemma to the case of $p(X)$ and $(a-X)p(X)$ by substituting X with aX , yielding the polynomials $p(aX)$ and $a \cdot (1-X) \cdot p(aX)$.

lemma *sign-changes-poly-times-root-minus-x*:
fixes $g :: 'a :: \text{linordered-idom poly}$ **and** $a :: 'a$


```

assumes  $nz: g \neq 0$  and  $pos: a > 0$ 
defines  $v \equiv \text{coeff-sign-changes}$ 
shows  $v ([:a, -1:] * g) - v g > 0 \wedge \text{odd} (v ([:a, -1:] * g) - v g)$ 
proof -
  have  $0 < v ([:1, -1:] * \text{reduce-root } a \ g) - v (\text{reduce-root } a \ g) \wedge$ 
     $\text{odd} (v ([:1, -1:] * \text{reduce-root } a \ g) - v (\text{reduce-root } a \ g))$ 
    using  $nz \ pos \ \text{unfolding } v\text{-def}$  by  $(\text{intro } \text{sign-changes-poly-times-one-minus-x})$ 
  simp-all
  also have  $v ([:1, -1:] * \text{reduce-root } a \ g) = v (\text{smult } a ([:1, -1:] * \text{reduce-root } a \ g))$ 
    unfolding  $v\text{-def}$  by  $(\text{simp add: } \text{coeff-sign-changes-smult } pos)$ 
  also have  $\text{smult } a ([:1, -1:] * \text{reduce-root } a \ g) = [:a:] * [:1, -1:] * \text{reduce-root } a \ g$ 
    by  $(\text{subst } \text{mult.assoc})$  simp
  also have  $[:a:] * [:1, -1:] = \text{reduce-root } a \ [:a, -1:]$ 
    by  $(\text{simp add: } \text{reduce-root-def } p\text{compose-pCons})$ 
  also have  $\dots * \text{reduce-root } a \ g = \text{reduce-root } a ([:a, -1:] * g)$ 
    unfolding  $\text{reduce-root-def}$  by  $(\text{simp only: } p\text{compose-mult})$ 
  also have  $v \dots = v ([:a, -1:] * g)$  by  $(\text{simp add: } v\text{-def } \text{coeff-sign-changes-reduce-root } pos)$ 
  also have  $v (\text{reduce-root } a \ g) = v g$  by  $(\text{simp add: } v\text{-def } \text{coeff-sign-changes-reduce-root } pos)$ 
  finally show  $?thesis$  .
qed

```

Finally, the difference of the number of coefficient sign changes and the number of positive roots is non-negative and even. This follows straightforwardly by induction over the roots.

lemma *descartes-sign-rule-aux*:

```

fixes  $p :: \text{real poly}$ 
assumes  $p \neq 0$ 
shows  $\text{coeff-sign-changes } p \geq \text{count-pos-roots } p \wedge$ 
   $\text{even} (\text{coeff-sign-changes } p - \text{count-pos-roots } p)$ 
using assms
proof  $(\text{induction } p \ \text{rule: } \text{poly-root-induct}[\text{where } P = \lambda a. a > 0])$ 
  case  $(\text{root } a \ p)$ 
    define  $q$  where  $q = [:a, -1:] * p$ 
    from  $\text{root.prem}$  have  $p: p \neq 0$  by auto
    with  $\text{root } p \ \text{sign-changes-poly-times-root-minus-x}[of \ p \ a]$ 
       $\text{count-roots-with-times-root}[of \ p \ \lambda x. x > 0 \ a]$  show  $?case$  by  $(\text{fold } q\text{-def})$ 
  fastforce
next
  case  $(\text{no-roots } p)$ 
    from  $\text{no-roots}$  have  $\text{pos-roots } p = \{\}$  by  $(\text{auto } \text{simp: } \text{roots-with-def})$ 
    hence  $[\text{simp}]: \text{count-pos-roots } p = 0$  by  $(\text{simp add: } \text{count-roots-with-def})$ 
    thus  $?case$  using  $\text{no-roots } \langle p \neq 0 \rangle \ \text{odd-coeff-sign-changes-imp-pos-roots}[of \ p]$ 
      by  $(\text{auto } \text{simp: } \text{roots-with-def})$ 
qed simp-all

```

The main theorem is then an obvious consequence

theorem *descartes-sign-rule*:

fixes $p :: \text{real poly}$

assumes $p \neq 0$

shows $\exists d. \text{even } d \wedge \text{coeff-sign-changes } p = \text{count-pos-roots } p + d$

proof

define d **where** $d = \text{coeff-sign-changes } p - \text{count-pos-roots } p$

show $\text{even } d \wedge \text{coeff-sign-changes } p = \text{count-pos-roots } p + d$

unfolding $d\text{-def}$ **using** *descartes-sign-rule-aux*[*OF* *assms*] **by** *auto*

qed

end

References

- [1] R. D. Arthan. Descartes' rule of signs by an easy induction. 2007.