

# Count the Number of Complex Roots

Wenda Li

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## Abstract

Based on evaluating Cauchy indices through remainder sequences [1] [2, Chapter 11], this entry provides an effective procedure to count the number of complex roots (with multiplicity) of a polynomial within a rectangle box or a half-plane. Potential applications of this entry include certified complex root isolation (of a polynomial) and testing the Routh-Hurwitz stability criterion (i.e., to check whether all the roots of some characteristic polynomial have negative real parts).

## 1 An alternative Sturm sequences

**theory** *Extended-Sturm* **imports**

*Sturm-Tarski.Sturm-Tarski*

*Winding-Number-Eval.Cauchy-Index-Theorem*

**begin**

The main purpose of this theory is to provide an effective way to compute  $\text{cindex} E a b f$  when  $f$  is a rational function. The idea is similar to and based on the evaluation of  $\text{cindex-poly}$  through  $[[?a < ?b; \text{poly } ?p \ ?a \neq 0; \text{poly } ?p \ ?b \neq 0]] \implies \text{cindex-poly } ?a \ ?b \ ?q \ ?p = \text{changes-itv-smods } ?a \ ?b \ ?p \ ?q$ .

This alternative version of remainder sequences is inspired by the paper "The Fundamental Theorem of Algebra made effective: an elementary real-algebraic proof via Sturm chains" by Michael Eisermann.

**hide-const** *Permutations.sign*

### 1.1 Misc

**lemma** *is-unit-pCons-ex-iff*:

**fixes**  $p::'a::\text{field poly}$

**shows**  $\text{is-unit } p \longleftrightarrow (\exists a. a \neq 0 \wedge p = [a:])$

*<proof>*

**lemma** *poly-gcd-iff*:

$\text{poly } (\text{gcd } p \ q) \ x=0 \longleftrightarrow \text{poly } p \ x=0 \wedge \text{poly } q \ x=0$

*<proof>*

**lemma** *eventually-poly-nz-at-within*:  
**fixes**  $x :: 'a::\{idom,euclidean-space\}$   
**assumes**  $p \neq 0$   
**shows** *eventually*  $(\lambda x. \text{poly } p \ x \neq 0)$  (at  $x$  within  $S$ )  
 $\langle$ *proof* $\rangle$

**lemma** *sgn-power*:  
**fixes**  $x::'a::linordered-idom$   
**shows**  $\text{sgn } (x^n) = (\text{if } n=0 \text{ then } 1 \text{ else if even } n \text{ then } |\text{sgn } x| \text{ else } \text{sgn } x)$   
 $\langle$ *proof* $\rangle$

**lemma** *poly-divide-filterlim-at-top*:  
**fixes**  $p \ q::\text{real poly}$   
**defines**  $ll \equiv (\text{if degree } q < \text{degree } p \text{ then}$   
 $\text{at } 0$   
 $\text{else if degree } q = \text{degree } p \text{ then}$   
 $\text{nhds } (\text{lead-coeff } q / \text{lead-coeff } p)$   
 $\text{else if sgn-pos-inf } q * \text{sgn-pos-inf } p > 0 \text{ then}$   
 $\text{at-top}$   
 $\text{else}$   
 $\text{at-bot})$   
**assumes**  $p \neq 0 \ q \neq 0$   
**shows** *filterlim*  $(\lambda x. \text{poly } q \ x / \text{poly } p \ x)$   $ll$  *at-top*  
 $\langle$ *proof* $\rangle$

**lemma** *poly-divide-filterlim-at-bot*:  
**fixes**  $p \ q::\text{real poly}$   
**defines**  $ll \equiv (\text{if degree } q < \text{degree } p \text{ then}$   
 $\text{at } 0$   
 $\text{else if degree } q = \text{degree } p \text{ then}$   
 $\text{nhds } (\text{lead-coeff } q / \text{lead-coeff } p)$   
 $\text{else if sgn-neg-inf } q * \text{sgn-neg-inf } p > 0 \text{ then}$   
 $\text{at-top}$   
 $\text{else}$   
 $\text{at-bot})$   
**assumes**  $p \neq 0 \ q \neq 0$   
**shows** *filterlim*  $(\lambda x. \text{poly } q \ x / \text{poly } p \ x)$   $ll$  *at-bot*  
 $\langle$ *proof* $\rangle$

## 1.2 Alternative definition of cross

**definition** *cross-alt* ::  $\text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{int}$  **where**  
 $\text{cross-alt } p \ q \ a \ b = |\text{sign } (\text{poly } p \ a) - \text{sign } (\text{poly } q \ a)| - |\text{sign } (\text{poly } p \ b) - \text{sign } (\text{poly } q \ b)|$

**lemma** *cross-alt-coprime-0*:  
**assumes**  $\text{coprime } p \ q \ p = 0 \vee q = 0$   
**shows**  $\text{cross-alt } p \ q \ a \ b = 0$   
 $\langle$ *proof* $\rangle$

**lemma** *cross-alt-0[simp]*:  $\text{cross-alt } 0 \ 0 \ a \ b = 0$   $\langle \text{proof} \rangle$

**lemma** *cross-alt-poly-commute*:  
 $\text{cross-alt } p \ q \ a \ b = \text{cross-alt } q \ p \ a \ b$   
 $\langle \text{proof} \rangle$

**lemma** *cross-alt-clear-n*:  
**assumes** *coprime*  $p \ q$   
**shows**  $\text{cross-alt } p \ q \ a \ b = \text{cross-alt } 1 \ (p * q) \ a \ b$   
 $\langle \text{proof} \rangle$

### 1.3 Alternative sign variation sequence

**fun** *changes-alt*::  $('a :: \text{linordered-idom}) \ \text{list} \Rightarrow \text{int}$  **where**  
 $\text{changes-alt } [] = 0$   
 $\text{changes-alt } [-] = 0$   
 $\text{changes-alt } (x1 \# x2 \# xs) = \text{abs}(\text{sign } x1 - \text{sign } x2) + \text{changes-alt } (x2 \# xs)$

**definition** *changes-alt-poly-at*::  $('a :: \text{linordered-idom}) \ \text{poly list} \Rightarrow 'a \Rightarrow \text{int}$  **where**  
 $\text{changes-alt-poly-at } ps \ a = \text{changes-alt } (\text{map } (\lambda p. \ \text{poly } p \ a) \ ps)$

**definition** *changes-alt-itv-smods*::  $\text{real} \Rightarrow \text{real} \Rightarrow \text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{int}$   
**where**  
 $\text{changes-alt-itv-smods } a \ b \ p \ q = (\text{let } ps = \text{smods } p \ q$   
 $\text{in } \text{changes-alt-poly-at } ps \ a - \text{changes-alt-poly-at } ps \ b)$

**lemma** *changes-alt-itv-smods-rec*:  
**assumes**  $a < b$  *coprime*  $p \ q$   
**shows**  $\text{changes-alt-itv-smods } a \ b \ p \ q = \text{cross-alt } p \ q \ a \ b + \text{changes-alt-itv-smods } a \ b \ q \ (- (p \ \text{mod } q))$   
 $\langle \text{proof} \rangle$

### 1.4 jumpF on polynomials

**definition** *jumpF-polyR*::  $\text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{real} \Rightarrow \text{real}$  **where**  
 $\text{jumpF-polyR } q \ p \ a = \text{jumpF } (\lambda x. \ \text{poly } q \ x / \text{poly } p \ x) \ (\text{at-right } a)$

**definition** *jumpF-polyL*::  $\text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{real} \Rightarrow \text{real}$  **where**  
 $\text{jumpF-polyL } q \ p \ a = \text{jumpF } (\lambda x. \ \text{poly } q \ x / \text{poly } p \ x) \ (\text{at-left } a)$

**definition** *jumpF-poly-top*::  $\text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{real}$  **where**  
 $\text{jumpF-poly-top } q \ p = \text{jumpF } (\lambda x. \ \text{poly } q \ x / \text{poly } p \ x) \ \text{at-top}$

**definition** *jumpF-poly-bot*::  $\text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{real}$  **where**  
 $\text{jumpF-poly-bot } q \ p = \text{jumpF } (\lambda x. \ \text{poly } q \ x / \text{poly } p \ x) \ \text{at-bot}$

**lemma** *jumpF-polyR-0[simp]*:  $\text{jumpF-polyR } 0 \ p \ a = 0$   $\text{jumpF-polyR } q \ 0 \ a = 0$   
 $\langle \text{proof} \rangle$



**lemma** *jumpF-polyL-smult-1*:

*jumpF-polyL (smult c q) p x = sgn c \* jumpF-polyL q p x*  
(proof)

**lemma** *jumpF-polyR-smult-1*:

*jumpF-polyR (smult c q) p x = sgn c \* jumpF-polyR q p x*  
(proof)

**lemma**

**shows** *jumpF-polyR-mod:jumpF-polyR q p x = jumpF-polyR (q mod p) p x* **and**  
*jumpF-polyL-mod:jumpF-polyL q p x = jumpF-polyL (q mod p) p x*  
(proof)

**lemma** *jumpF-poly-top-0[simp]*: *jumpF-poly-top 0 p = 0 jumpF-poly-top q 0 = 0*  
(proof)

**lemma** *jumpF-poly-bot-0[simp]*: *jumpF-poly-bot 0 p = 0 jumpF-poly-bot q 0 = 0*  
(proof)

**lemma** *jumpF-poly-top-code*:

*jumpF-poly-top q p = (if p≠0 ∧ q≠0 ∧ degree q > degree p then*  
*if sgn-pos-inf q \* sgn-pos-inf p > 0 then 1/2 else -1/2 else 0)*  
(proof)

**lemma** *jumpF-poly-bot-code*:

*jumpF-poly-bot q p = (if p≠0 ∧ q≠0 ∧ degree q > degree p then*  
*if sgn-neg-inf q \* sgn-neg-inf p > 0 then 1/2 else -1/2 else 0)*  
(proof)

## 1.5 The extended Cauchy index on polynomials

**definition** *cindex-polyE*:: *real ⇒ real ⇒ real poly ⇒ real poly ⇒ real* **where**

*cindex-polyE a b q p = jumpF-polyR q p a + cindex-poly a b q p - jumpF-polyL*  
*q p b*

**definition** *cindex-poly-ubd*:: *real poly ⇒ real poly ⇒ int* **where**

*cindex-poly-ubd q p = (THE l. (∀<sub>F</sub> r in at-top. cindexE (-r) r (λx. poly q x / poly*  
*p x) = of-int l))*

**lemma** *cindex-polyE-0[simp]*: *cindex-polyE a b 0 p = 0 cindex-polyE a b q 0 = 0*  
(proof)

**lemma** *cindex-polyE-mult-cancel*:

**fixes** *p q p'*:: *real poly*  
**assumes** *p' ≠ 0*

**shows**  $cindex\text{-}polyE\ a\ b\ (p' * q)\ (p' * p) = cindex\text{-}polyE\ a\ b\ q\ p$   
 ⟨proof⟩

**lemma**  $cindexE\text{-}eq\text{-}cindex\text{-}polyE$ :

**assumes**  $a < b$

**shows**  $cindexE\ a\ b\ (\lambda x. poly\ q\ x / poly\ p\ x) = cindex\text{-}polyE\ a\ b\ q\ p$   
 ⟨proof⟩

**lemma**  $cindex\text{-}polyE\text{-}cross$ :

**fixes**  $p::real\ poly$  **and**  $a\ b::real$

**assumes**  $a < b$

**shows**  $cindex\text{-}polyE\ a\ b\ 1\ p = cross\text{-}alt\ 1\ p\ a\ b / 2$   
 ⟨proof⟩

**lemma**  $cindex\text{-}polyE\text{-}inverse\text{-}add$ :

**fixes**  $p\ q::real\ poly$

**assumes**  $cp::coprime\ p\ q$

**shows**  $cindex\text{-}polyE\ a\ b\ q\ p + cindex\text{-}polyE\ a\ b\ p\ q = cindex\text{-}polyE\ a\ b\ 1\ (q*p)$   
 ⟨proof⟩

**lemma**  $cindex\text{-}polyE\text{-}inverse\text{-}add\text{-}cross$ :

**fixes**  $p\ q::real\ poly$

**assumes**  $a < b\ coprime\ p\ q$

**shows**  $cindex\text{-}polyE\ a\ b\ q\ p + cindex\text{-}polyE\ a\ b\ p\ q = cross\text{-}alt\ p\ q\ a\ b / 2$   
 ⟨proof⟩

**lemma**  $cindex\text{-}polyE\text{-}smult\text{-}1$ :

**fixes**  $p\ q::real\ poly$  **and**  $c::real$

**shows**  $cindex\text{-}polyE\ a\ b\ (smult\ c\ q)\ p = (sgn\ c) * cindex\text{-}polyE\ a\ b\ q\ p$   
 ⟨proof⟩

**lemma**  $cindex\text{-}polyE\text{-}mod$ :

**fixes**  $p\ q::real\ poly$

**shows**  $cindex\text{-}polyE\ a\ b\ q\ p = cindex\text{-}polyE\ a\ b\ (q\ mod\ p)\ p$   
 ⟨proof⟩

**lemma**  $cindex\text{-}polyE\text{-}rec$ :

**fixes**  $p\ q::real\ poly$

**assumes**  $a < b\ coprime\ p\ q$

**shows**  $cindex\text{-}polyE\ a\ b\ q\ p = cross\text{-}alt\ q\ p\ a\ b / 2 + cindex\text{-}polyE\ a\ b\ (- (p\ mod\ q))\ q$   
 ⟨proof⟩

**lemma**  $cindex\text{-}polyE\text{-}changes\text{-}alt\text{-}itv\text{-}mods$ :

**assumes**  $a < b\ coprime\ p\ q$

**shows**  $cindex\text{-}polyE\ a\ b\ q\ p = changes\text{-}alt\text{-}itv\text{-}smods\ a\ b\ p\ q / 2$  ⟨proof⟩

**lemma**  $cindex\text{-}poly\text{-}ubd\text{-}eventually$ :

**shows**  $\forall_F\ r\ in\ at\text{-}top. cindexE\ (-r)\ r\ (\lambda x. poly\ q\ x / poly\ p\ x) = of\text{-}int\ (cindex\text{-}poly\text{-}ubd$

$q \ p)$   
 $\langle proof \rangle$

**lemma** *cindex-poly-ubd-0*:  
  **assumes**  $p=0 \vee q=0$   
  **shows**  $cindex\text{-}poly\text{-}ubd \ q \ p = 0$   
 $\langle proof \rangle$

**lemma** *cindex-poly-ubd-code*:  
  **shows**  $cindex\text{-}poly\text{-}ubd \ q \ p = changes\text{-}R\text{-}smods \ p \ q$   
 $\langle proof \rangle$

**lemma** *cindexE-ubd-poly*:  $cindexE\text{-}ubd \ (\lambda x. \ poly \ q \ x / poly \ p \ x) = cindex\text{-}poly\text{-}ubd$   
 $q \ p$   
 $\langle proof \rangle$

**end**

## 2 More useful lemmas related polynomials

**theory** *More-Polynomials imports*  
  *Winding-Number-Eval.Missing-Algebraic*  
  *Winding-Number-Eval.Missing-Transcendental*  
  *Sturm-Tarski.PolyMisc*  
  *Budan-Fourier.BF-Misc*  
**begin**

### 2.1 More about order

**lemma** *order-normalize[simp]*:  $order \ x \ (normalize \ p) = order \ x \ p$   
 $\langle proof \rangle$

**lemma** *order-gcd*:  
  **assumes**  $p \neq 0 \ q \neq 0$   
  **shows**  $order \ x \ (gcd \ p \ q) = min \ (order \ x \ p) \ (order \ x \ q)$   
 $\langle proof \rangle$

**lemma** *pderiv-power*:  $pderiv \ (p \ ^ \ n) = smult \ (of\text{-}nat \ n) \ (p \ ^ \ (n-1)) * pderiv \ p$   
 $\langle proof \rangle$

**lemma** *order-pderiv*:  
  **fixes**  $p::'a::\{idom,semiring\text{-}char\text{-}0\} \ poly$   
  **assumes**  $p \neq 0 \ poly \ p \ x = 0$   
  **shows**  $order \ x \ p = Suc \ (order \ x \ (pderiv \ p)) \ \langle proof \rangle$

## 2.2 More about *rsquarefree*

**lemma** *rsquarefree-0[simp]*:  $\neg$  *rsquarefree* 0  
⟨*proof*⟩

**lemma** *rsquarefree-times*:  
 **assumes** *rsquarefree* (p\*q)  
 **shows** *rsquarefree* q ⟨*proof*⟩

**lemma** *rsquarefree-smult-iff*:  
 **assumes**  $s \neq 0$   
 **shows** *rsquarefree* (smult s p)  $\longleftrightarrow$  *rsquarefree* p  
 ⟨*proof*⟩

**lemma** *card-proots-within-rsquarefree*:  
 **assumes** *rsquarefree* p  
 **shows** *proots-count* p s = *card* (*proots-within* p s) ⟨*proof*⟩

**lemma** *rsquarefree-gcd-pderiv*:  
 **fixes** p::'a::{*factorial-ring-gcd*,*semiring-char-0*} *poly*  
 **assumes**  $p \neq 0$   
 **shows** *rsquarefree* (p div (gcd p (pderiv p)))  
 ⟨*proof*⟩

**lemma** *poly-gcd-pderiv-iff*:  
 **fixes** p::'a::{*semiring-char-0*,*factorial-ring-gcd*} *poly*  
 **shows** *poly* (p div (gcd p (pderiv p)))  $x=0 \longleftrightarrow$  *poly* p  $x=0$   
 ⟨*proof*⟩

## 2.3 Composition of a polynomial and a circular path

**lemma** *poly-circlepath-tan-eq*:  
 **fixes** z0::*complex* **and** r::*real* **and** p::*complex poly*  
 **defines** q1  $\equiv$  *fcompose* p [:(z0+r)\*i,z0-r:] [i,1:] **and** q2  $\equiv$  [i,1:] ^ *degree* p  
 **assumes**  $0 \leq t \leq 1$   $t \neq 1/2$   
 **shows** *poly* p (*circlepath* z0 r t) = *poly* q1 (*tan* (pi\*t)) / *poly* q2 (*tan* (pi\*t))  
 (is ?L = ?R)  
 ⟨*proof*⟩

**end**

## 3 Procedures to count the number of complex roots

**theory** *Count-Complex-Roots* **imports**  
 *Winding-Number-Eval*.*Winding-Number-Eval*  
 *Extended-Sturm*  
 *More-Polynomials*  
 *Budan-Fourier*.*Sturm-Multiple-Roots*  
**begin**



### 3.1 Misc

**corollary** *path-image-part-circlepath-subset*:

**assumes**  $r \geq 0$   
**shows**  $\text{path-image}(\text{part-circlepath } z \ r \ st \ tt) \subseteq \text{sphere } z \ r$   
(*proof*)

**proposition** *in-path-image-part-circlepath*:

**assumes**  $w \in \text{path-image}(\text{part-circlepath } z \ r \ st \ tt)$   $0 \leq r$   
**shows**  $\text{norm}(w - z) = r$   
(*proof*)

**lemma** *infinite-ball*:

**fixes**  $a :: 'a::\text{euclidean-space}$   
**assumes**  $r > 0$   
**shows** *infinite* ( $\text{ball } a \ r$ )  
(*proof*)

**lemma** *infinite-cball*:

**fixes**  $a :: 'a::\text{euclidean-space}$   
**assumes**  $r > 0$   
**shows** *infinite* ( $\text{cball } a \ r$ )  
(*proof*)

**lemma** *infinite-sphere*:

**fixes**  $a :: \text{complex}$   
**assumes**  $r > 0$   
**shows** *infinite* ( $\text{sphere } a \ r$ )  
(*proof*)

**lemma** *infinite-halfspace-Im-gt*: *infinite*  $\{x. \text{Im } x > b\}$

(*proof*)

**lemma** (*in ring-1*) *Ints-minus2*:  $- a \in \mathbb{Z} \implies a \in \mathbb{Z}$

(*proof*)

**lemma** *dvd-divide-Ints-iff*:

$b \text{ dvd } a \vee b=0 \iff \text{of-int } a / \text{of-int } b \in (\mathbb{Z} :: 'a :: \{\text{field,ring-char-0}\} \text{ set})$   
(*proof*)

**lemma** *of-int-div-field*:

**assumes**  $d \text{ dvd } n$   
**shows** ( $\text{of-int}::\implies 'a::\text{field-char-0}$ )  $(n \text{ div } d) = \text{of-int } n / \text{of-int } d$   
(*proof*)

**lemma** *powr-eq-1-iff*:

**assumes**  $a > 0$   
**shows** ( $a::\text{real}$ )  $\text{powr } b = 1 \iff a=1 \vee b=0$

*<proof>*

**lemma** *tan-inj-pi*:

$-(\pi/2) < x \implies x < \pi/2 \implies -(\pi/2) < y \implies y < \pi/2 \implies \tan x = \tan y \implies x = y$

*<proof>*

**lemma** *finite-ReZ-segments-poly-circlepath*:

*finite-ReZ-segments (poly p o circlepath z0 r) 0*

*<proof>*

### 3.2 Some useful conformal/*bij-betw* properties

**lemma** *bij-betw-plane-ball*:*bij-betw*  $(\lambda x. (i-x)/(i+x)) \{x. \text{Im } x > 0\}$  (ball 0 1)

*<proof>*

**lemma** *bij-betw-axis-sphere*:*bij-betw*  $(\lambda x. (i-x)/(i+x)) \{x. \text{Im } x = 0\}$  (sphere 0 1 - { -1 })

*<proof>*

**lemma** *bij-betw-ball-uball*:

**assumes**  $r > 0$

**shows** *bij-betw*  $(\lambda x. \text{complex-of-real } r*x + z0)$  (ball 0 1) (ball z0 r)

*<proof>*

**lemma** *bij-betw-sphere-usphere*:

**assumes**  $r > 0$

**shows** *bij-betw*  $(\lambda x. \text{complex-of-real } r*x + z0)$  (sphere 0 1) (sphere z0 r)

*<proof>*

**lemma** *proots-ball-plane-eq*:

**defines**  $q1 \equiv [i, -1:]$  **and**  $q2 \equiv [i, 1:]$

**assumes**  $p \neq 0$

**shows** *proots-count*  $p$  (ball 0 1) = *proots-count* (fcompose  $p$   $q1$   $q2$ )  $\{x. 0 < \text{Im } x\}$

*<proof>*

**lemma** *proots-sphere-axis-eq*:

**defines**  $q1 \equiv [i, -1:]$  **and**  $q2 \equiv [i, 1:]$

**assumes**  $p \neq 0$

**shows** *proots-count*  $p$  (sphere 0 1 - { -1 }) = *proots-count* (fcompose  $p$   $q1$   $q2$ )  $\{x. 0 = \text{Im } x\}$

*<proof>*

**lemma** *proots-card-ball-plane-eq*:

**defines**  $q1 \equiv [i, -1:]$  **and**  $q2 \equiv [i, 1:]$

**assumes**  $p \neq 0$

**shows** *card* (*proots-within*  $p$  (ball 0 1)) = *card* (*proots-within* (fcompose  $p$   $q1$   $q2$ ))

$\{x. 0 < \text{Im } x\}$   
 $\langle \text{proof} \rangle$

**lemma** *proots-card-sphere-axis-eq*:  
**defines**  $q1 \equiv [i, -1:]$  **and**  $q2 \equiv [i, 1:]$   
**assumes**  $p \neq 0$   
**shows**  $\text{card } (\text{proots-within } p \ (\text{sphere } 0 \ 1 \ - \ \{-1\}))$   
 $= \text{card } (\text{proots-within } (\text{fcompose } p \ q1 \ q2) \ \{x. 0 = \text{Im } x\})$   
 $\langle \text{proof} \rangle$

**lemma** *proots-uball-eq*:  
**fixes**  $z0::\text{complex}$  **and**  $r::\text{real}$   
**defines**  $q \equiv [z0, \text{of-real } r:]$   
**assumes**  $p \neq 0$  **and**  $r > 0$   
**shows**  $\text{proots-count } p \ (\text{ball } z0 \ r) = \text{proots-count } (p \circ_p \ q) \ (\text{ball } 0 \ 1)$   
 $\langle \text{proof} \rangle$

**lemma** *proots-card-uball-eq*:  
**fixes**  $z0::\text{complex}$  **and**  $r::\text{real}$   
**defines**  $q \equiv [z0, \text{of-real } r:]$   
**assumes**  $r > 0$   
**shows**  $\text{card } (\text{proots-within } p \ (\text{ball } z0 \ r)) = \text{card } (\text{proots-within } (p \circ_p \ q) \ (\text{ball } 0 \ 1))$   
 $\langle \text{proof} \rangle$

**lemma** *proots-card-usphere-eq*:  
**fixes**  $z0::\text{complex}$  **and**  $r::\text{real}$   
**defines**  $q \equiv [z0, \text{of-real } r:]$   
**assumes**  $r > 0$   
**shows**  $\text{card } (\text{proots-within } p \ (\text{sphere } z0 \ r)) = \text{card } (\text{proots-within } (p \circ_p \ q) \ (\text{sphere } 0 \ 1))$   
 $\langle \text{proof} \rangle$

### 3.3 Combining two real polynomials into a complex one

**definition** *cpoly-of*::  $\text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{complex poly}$  **where**  
 $\text{cpoly-of } pR \ pI = \text{map-poly of-real } pR + \text{smult } i \ (\text{map-poly of-real } pI)$

**lemma** *cpoly-of-eq-0-iff* [iff]:  
 $\text{cpoly-of } pR \ pI = 0 \iff pR = 0 \wedge pI = 0$   
 $\langle \text{proof} \rangle$

**lemma** *cpoly-of-decompose*:  
 $p = \text{cpoly-of } (\text{map-poly } \text{Re } p) \ (\text{map-poly } \text{Im } p)$   
 $\langle \text{proof} \rangle$

**lemma** *cpoly-of-dist-right*:  
 $\text{cpoly-of } (pR * g) \ (pI * g) = \text{cpoly-of } pR \ pI * (\text{map-poly of-real } g)$   
 $\langle \text{proof} \rangle$

**lemma** *poly-cpoly-of-real*:

$\text{poly } (\text{cpoly-of } pR \ pI) \ (\text{of-real } x) = \text{Complex } (\text{poly } pR \ x) \ (\text{poly } pI \ x)$   
 ⟨proof⟩

**lemma** *poly-cpoly-of-real-iff*:

**shows**  $\text{poly } (\text{cpoly-of } pR \ pI) \ (\text{of-real } t) = 0 \iff \text{poly } pR \ t = 0 \wedge \text{poly } pI \ t = 0$   
 ⟨proof⟩

**lemma** *order-cpoly-gcd-eq*:

**assumes**  $pR \neq 0 \vee pI \neq 0$   
**shows**  $\text{order } t \ (\text{cpoly-of } pR \ pI) = \text{order } t \ (\text{gcd } pR \ pI)$   
 ⟨proof⟩

### 3.4 Number of roots on a (bounded or unbounded) segment

— 1 dimensional hyperplane

**definition** *unbounded-line*:: 'a::real-vector  $\Rightarrow$  'a  $\Rightarrow$  'a set **where**

$\text{unbounded-line } a \ b = \{\{x. \exists u::\text{real}. x = (1 - u) *_R a + u *_R b\}\}$

**definition** *roots-line-card*:: complex poly  $\Rightarrow$  complex  $\Rightarrow$  complex  $\Rightarrow$  nat **where**

$\text{roots-line-card } p \ st \ tt = \text{card } (\text{roots-within } p \ (\text{open-segment } st \ tt))$

**definition** *roots-unbounded-line-card*:: complex poly  $\Rightarrow$  complex  $\Rightarrow$  complex  $\Rightarrow$  nat **where**

$\text{roots-unbounded-line-card } p \ st \ tt = \text{card } (\text{roots-within } p \ (\text{unbounded-line } st \ tt))$

**definition** *roots-unbounded-line* :: complex poly  $\Rightarrow$  complex  $\Rightarrow$  complex  $\Rightarrow$  nat **where**

$\text{roots-unbounded-line } p \ st \ tt = \text{roots-count } p \ (\text{unbounded-line } st \ tt)$

**lemma** *card-proots-open-segments*:

**assumes**  $\text{poly } p \ st \neq 0 \ \text{poly } p \ tt \neq 0$

**shows**  $\text{card } (\text{roots-within } p \ (\text{open-segment } st \ tt)) =$

$(\text{let } pc = \text{pcompose } p \ [ :st, tt - st:];$

$pR = \text{map-poly } \text{Re } pc;$

$pI = \text{map-poly } \text{Im } pc;$

$g = \text{gcd } pR \ pI$

$\text{in changes-itu-smods } 0 \ 1 \ g \ (\text{pderiv } g)) \ (\text{is } ?L = ?R)$

⟨proof⟩

**lemma** *unbounded-line-closed-segment*: closed-segment a b  $\subseteq$  unbounded-line a b

⟨proof⟩

**lemma** *card-proots-unbounded-line*:

**assumes**  $st \neq tt$

**shows**  $\text{card } (\text{roots-within } p \ (\text{unbounded-line } st \ tt)) =$

$(\text{let } pc = \text{pcompose } p \ [ :st, tt - st:];$

$pR = \text{map-poly } \text{Re } pc;$

$pI = \text{map-poly } \text{Im } pc;$   
 $g = \text{gcd } pR \ pI$   
 $\text{in nat } (\text{changes-R-smods } g \ (\text{pderiv } g)) \ (\text{is } ?L = ?R)$

*<proof>*

**lemma** *roots-unbounded-line:*

**assumes**  $st \neq tt \ p \neq 0$

**shows**  $(\text{roots-count } p \ (\text{unbounded-line } st \ tt)) =$   
 $(\text{let } pc = \text{pcompose } p \ [ :st, tt - st:];$

$pR = \text{map-poly } \text{Re } pc;$

$pI = \text{map-poly } \text{Im } pc;$

$g = \text{gcd } pR \ pI$

$\text{in nat } (\text{changes-R-smods-ext } g \ (\text{pderiv } g)) \ (\text{is } ?L = ?R)$

*<proof>*

**lemma** *roots-unbounded-line-card-code*[code]:

*roots-unbounded-line-card*  $p \ st \ tt =$

*(if*  $st \neq tt$  *then*

$(\text{let } pc = \text{pcompose } p \ [ :st, tt - st:];$

$pR = \text{map-poly } \text{Re } pc;$

$pI = \text{map-poly } \text{Im } pc;$

$g = \text{gcd } pR \ pI$

$\text{in nat } (\text{changes-R-smods } g \ (\text{pderiv } g))$

*else*

$\text{Code.abort } (\text{STR } \text{"roots-unbounded-line-card fails due to invalid$

$\text{hyperplanes."})$

$(\lambda-. \text{roots-unbounded-line-card } p \ st \ tt))$

*<proof>*

**lemma** *roots-unbounded-line-code*[code]:

*roots-unbounded-line*  $p \ st \ tt =$

*( if*  $st \neq tt$  *then*

*if*  $p \neq 0$  *then*

$(\text{let } pc = \text{pcompose } p \ [ :st, tt - st:];$

$pR = \text{map-poly } \text{Re } pc;$

$pI = \text{map-poly } \text{Im } pc;$

$g = \text{gcd } pR \ pI$

$\text{in nat } (\text{changes-R-smods-ext } g \ (\text{pderiv } g))$

*else*

$\text{Code.abort } (\text{STR } \text{"roots-unbounded-line fails due to } p=0\text{"})$

$(\lambda-. \text{roots-unbounded-line } p \ st \ tt)$

*else*

$\text{Code.abort } (\text{STR } \text{"roots-unbounded-line fails due to invalid$

$\text{hyperplanes."})$

$(\lambda-. \text{roots-unbounded-line } p \ st \ tt) )$

*<proof>*

### 3.5 Checking if there a polynomial root on a closed segment

**definition** *no-roots-line* :: *complex poly*  $\Rightarrow$  *complex*  $\Rightarrow$  *complex*  $\Rightarrow$  *bool* **where**  
*no-roots-line* *p st tt* = (*roots-within* *p* (*closed-segment* *st tt*) = {})

**lemma** *no-roots-line-code*[*code*]: *no-roots-line* *p st tt* = (*if* *poly* *p st*  $\neq 0$   $\wedge$  *poly* *p tt*  $\neq 0$  *then*

(*let* *pc* = *pcompose* *p* [:*st*, *tt* - *st*:];  
*pR* = *map-poly* *Re* *pc*;  
*pI* = *map-poly* *Im* *pc*;  
*g* = *gcd* *pR* *pI*  
*in if* *changes-itv-smods* 0 1 *g* (*pderiv* *g*) = 0 *then* *True* *else* *False*)

*else* *False*)

(*is* ?*L* = ?*R*)

*<proof>*

### 3.6 Counting roots in a rectangle

**definition** *roots-rectangle* :: *complex poly*  $\Rightarrow$  *complex*  $\Rightarrow$  *complex*  $\Rightarrow$  *nat* **where**  
*roots-rectangle* *p lb ub* = *roots-count* *p* (*box* *lb ub*)

**lemma** *closed-segment-imp-Re-Im*:

**fixes** *x* :: *complex*

**assumes** *x*  $\in$  *closed-segment* *lb ub*

**shows** *Re* *lb*  $\leq$  *Re* *ub*  $\implies$  *Re* *lb*  $\leq$  *Re* *x*  $\wedge$  *Re* *x*  $\leq$  *Re* *ub*

*Im* *lb*  $\leq$  *Im* *ub*  $\implies$  *Im* *lb*  $\leq$  *Im* *x*  $\wedge$  *Im* *x*  $\leq$  *Im* *ub*

*<proof>*

**lemma** *closed-segment-degen-complex*:

[[*Re* *lb* = *Re* *ub*; *Im* *lb*  $\leq$  *Im* *ub* ]]

$\implies$  *x*  $\in$  *closed-segment* *lb ub*  $\iff$  *Re* *x* = *Re* *lb*  $\wedge$  *Im* *lb*  $\leq$  *Im* *x*  $\wedge$  *Im* *x*  $\leq$  *Im* *ub*

*ub*

[[*Im* *lb* = *Im* *ub*; *Re* *lb*  $\leq$  *Re* *ub* ]]

$\implies$  *x*  $\in$  *closed-segment* *lb ub*  $\iff$  *Im* *x* = *Im* *lb*  $\wedge$  *Re* *lb*  $\leq$  *Re* *x*  $\wedge$  *Re* *x*  $\leq$  *Re* *ub*

*ub*

*<proof>*

**lemma** *complex-box-ne-empty*:

**fixes** *a b* :: *complex*

**shows**

*cbox* *a b*  $\neq$  {}  $\iff$  (*Re* *a*  $\leq$  *Re* *b*  $\wedge$  *Im* *a*  $\leq$  *Im* *b*)

*box* *a b*  $\neq$  {}  $\iff$  (*Re* *a*  $<$  *Re* *b*  $\wedge$  *Im* *a*  $<$  *Im* *b*)

*<proof>*

**lemma** *roots-rectangle-code1*:

*roots-rectangle* *p lb ub* = (*if* *Re* *lb*  $<$  *Re* *ub*  $\wedge$  *Im* *lb*  $<$  *Im* *ub* *then*

*if* *p*  $\neq 0$  *then*

*if* *no-roots-line* *p lb* (*Complex* (*Re* *ub*) (*Im* *lb*))

$\wedge$  *no-roots-line* *p* (*Complex* (*Re* *ub*) (*Im* *lb*)) *ub*

```

    ∧ no-roots-line p ub (Complex (Re lb) (Im ub))
    ∧ no-roots-line p (Complex (Re lb) (Im ub)) lb then
    (
    let p1 = pcompose p [:lb, Complex (Re ub - Re lb) 0:];
        pR1 = map-poly Re p1; pI1 = map-poly Im p1; gc1 = gcd pR1 pI1;
        p2 = pcompose p [:Complex (Re ub) (Im lb), Complex 0 (Im ub - Im
lb)];];
        pR2 = map-poly Re p2; pI2 = map-poly Im p2; gc2 = gcd pR2 pI2;
        p3 = pcompose p [:ub, Complex (Re lb - Re ub) 0:];
        pR3 = map-poly Re p3; pI3 = map-poly Im p3; gc3 = gcd pR3 pI3;
        p4 = pcompose p [:Complex (Re lb) (Im ub), Complex 0 (Im lb - Im
ub)];];
        pR4 = map-poly Re p4; pI4 = map-poly Im p4; gc4 = gcd pR4 pI4
in
    nat ( - (changes-alt-itv-smods 0 1 (pR1 div gc1) (pI1 div gc1)
    + changes-alt-itv-smods 0 1 (pR2 div gc2) (pI2 div gc2)
    + changes-alt-itv-smods 0 1 (pR3 div gc3) (pI3 div gc3)
    + changes-alt-itv-smods 0 1 (pR4 div gc4) (pI4 div gc4)) div 4)
    )
    else Code.abort (STR "proots-rectangle fails when there is a root on the
border.")
    (λ-. proots-rectangle p lb ub)
    else Code.abort (STR "proots-rectangle fails when p=0.")
    (λ-. proots-rectangle p lb ub)
    else 0)
⟨proof⟩

```

**lemma** *proots-rectangle-code2*[code]:

```

    proots-rectangle p lb ub = (if Re lb < Re ub ∧ Im lb < Im ub then
    if p≠0 then
    if poly p lb ≠ 0 ∧ poly p (Complex (Re ub) (Im lb)) ≠ 0
    ∧ poly p ub ≠ 0 ∧ poly p (Complex (Re lb) (Im ub)) ≠ 0
    then
    (let p1 = pcompose p [:lb, Complex (Re ub - Re lb) 0:];
        pR1 = map-poly Re p1; pI1 = map-poly Im p1; gc1 = gcd pR1 pI1;
        p2 = pcompose p [:Complex (Re ub) (Im lb), Complex 0 (Im ub - Im
lb)];];
        pR2 = map-poly Re p2; pI2 = map-poly Im p2; gc2 = gcd pR2 pI2;
        p3 = pcompose p [:ub, Complex (Re lb - Re ub) 0:];
        pR3 = map-poly Re p3; pI3 = map-poly Im p3; gc3 = gcd pR3 pI3;
        p4 = pcompose p [:Complex (Re lb) (Im ub), Complex 0 (Im lb - Im
ub)];];
        pR4 = map-poly Re p4; pI4 = map-poly Im p4; gc4 = gcd pR4 pI4
in
    if changes-itv-smods 0 1 gc1 (pderiv gc1) = 0
    ∧ changes-itv-smods 0 1 gc2 (pderiv gc2) = 0
    ∧ changes-itv-smods 0 1 gc3 (pderiv gc3) = 0
    ∧ changes-itv-smods 0 1 gc4 (pderiv gc4) = 0

```

```

then
  nat ( - (changes-alt-itv-smods 0 1 (pR1 div gc1) (pI1 div gc1)
    + changes-alt-itv-smods 0 1 (pR2 div gc2) (pI2 div gc2)
    + changes-alt-itv-smods 0 1 (pR3 div gc3) (pI3 div gc3)
    + changes-alt-itv-smods 0 1 (pR4 div gc4) (pI4 div gc4)) div 4)
  else Code.abort (STR "proots-rectangle fails when there is a root on
the border.")
      (λ-. proots-rectangle p lb ub)
  else Code.abort (STR "proots-rectangle fails when there is a root on the
border.")
      (λ-. proots-rectangle p lb ub)
  else Code.abort (STR "proots-rectangle fails when p=0.")
      (λ-. proots-rectangle p lb ub)
  else 0)
⟨proof⟩

```

### 3.7 Polynomial roots on the upper half-plane

— Roots counted WITH multiplicity

**definition** *proots-upper* :: *complex poly* ⇒ *nat* **where**  
*proots-upper* p = *proots-count* p {z. *Im* z > 0}

— Roots counted WITHOUT multiplicity

**definition** *proots-upper-card* :: *complex poly* ⇒ *nat* **where**  
*proots-upper-card* p = *card* (*proots-within* p {x. *Im* x > 0})

**lemma** *Im-Ln-tendsto-at-top*: ((λx. *Im* (Ln (Complex a x))) → pi/2 ) *at-top*  
⟨proof⟩

**lemma** *Im-Ln-tendsto-at-bot*: ((λx. *Im* (Ln (Complex a x))) → - pi/2 ) *at-bot*

⟨proof⟩

**lemma** *Re-winding-number-tendsto-part-circlepath*:

**shows** ((λr. *Re* (*winding-number* (*part-circlepath* z0 r 0 pi ) a)) → 1/2 )  
*at-top*  
⟨proof⟩

**lemma** *not-image-at-top-poly-part-circlepath*:

**assumes** *degree* p > 0  
**shows** ∀<sub>F</sub> r *in* *at-top*. b ∉ *path-image* (*poly* p o *part-circlepath* z0 r st tt)  
⟨proof⟩

**lemma** *not-image-poly-part-circlepath*:

**assumes** *degree* p > 0  
**shows** ∃ r > 0. b ∉ *path-image* (*poly* p o *part-circlepath* z0 r st tt)  
⟨proof⟩

**lemma** *Re-winding-number-poly-part-circlepath*:



**assumes**  $\text{degree } p > 0$   
**shows**  $((\lambda r. \text{Re } (\text{winding-number } (\text{poly } p \text{ o part-circlepath } z0 \ r \ 0 \ \pi) \ 0)) \longrightarrow \text{degree } p / 2)$  *at-top*  
 $\langle \text{proof} \rangle$

**lemma** *Re-winding-number-poly-linepth*:  
**fixes**  $pp::\text{complex poly}$   
**defines**  $g \equiv (\lambda r. \text{poly } pp \text{ o linepath } (-r) \ (\text{of-real } r))$   
**assumes**  $\text{lead-coeff } pp = 1$  **and**  $\text{no-real-zero}:\forall x \in \text{roots } pp. \text{Im } x \neq 0$   
**shows**  $((\lambda r. 2 * \text{Re } (\text{winding-number } (g \ r) \ 0) + \text{cindex-pathE } (g \ r) \ 0) \longrightarrow 0)$  *at-top*  
 $\langle \text{proof} \rangle$

**lemma** *roots-upper-cindex-eq*:  
**assumes**  $\text{lead-coeff } p = 1$  **and**  $\text{no-real-roots}:\forall x \in \text{roots } p. \text{Im } x \neq 0$   
**shows**  $\text{roots-upper } p = (\text{degree } p - \text{cindex-poly-ubd } (\text{map-poly } \text{Im } p) (\text{map-poly } \text{Re } p)) / 2$   
 $\langle \text{proof} \rangle$

**lemma** *cindexE-roots-on-horizontal-border*:  
**fixes**  $a::\text{complex and } s::\text{real}$   
**defines**  $g \equiv \text{linepath } a \ (a + \text{of-real } s)$   
**assumes**  $pqr:p = q * r$  **and**  $r\text{-monic}:\text{lead-coeff } r = 1$  **and**  $r\text{-roots}:\forall x \in \text{roots } r. \text{Im } x = \text{Im } a$   
**shows**  $\text{cindexE } lb \ ub \ (\lambda t. \text{Im } ((\text{poly } p \text{ o } g) \ t) / \text{Re } ((\text{poly } p \text{ o } g) \ t)) = \text{cindexE } lb \ ub \ (\lambda t. \text{Im } ((\text{poly } q \text{ o } g) \ t) / \text{Re } ((\text{poly } q \text{ o } g) \ t))$   
 $\langle \text{proof} \rangle$

**lemma** *poly-decompose-by-roots*:  
**fixes**  $p :: 'a::\text{idom poly}$   
**assumes**  $p \neq 0$   
**shows**  $\exists q \ r. p = q * r \wedge \text{lead-coeff } q = 1 \wedge (\forall x \in \text{roots } q. P \ x) \wedge (\forall x \in \text{roots } r. \neg P \ x)$   $\langle \text{proof} \rangle$

**lemma** *roots-upper-cindex-eq'*:  
**assumes**  $\text{lead-coeff } p = 1$   
**shows**  $\text{roots-upper } p = (\text{degree } p - \text{roots-count } p \ \{x. \text{Im } x = 0\} - \text{cindex-poly-ubd } (\text{map-poly } \text{Im } p) (\text{map-poly } \text{Re } p)) / 2$   
 $\langle \text{proof} \rangle$

**lemma** *roots-within-upper-squarefree*:  
**assumes**  $r\text{squarefree } p$   
**shows**  $\text{card } (\text{roots-within } p \ \{x. \text{Im } x > 0\}) = (\text{let } pp = \text{smult } (\text{inverse } (\text{lead-coeff } p)) \ p; pI = \text{map-poly } \text{Im } pp; pR = \text{map-poly } \text{Re } pp;$

```

      g = gcd pR pI
    in
      nat ((degree p - changes-R-smods g (pderiv g) - changes-R-smods pR
pI) div 2)
    )
  ⟨proof⟩

```

**lemma** *proots-upper-code1*[code]:

```

proots-upper p =
  (if p ≠ 0 then
    (let pp=smult (inverse (lead-coeff p)) p;
      pI=map-poly Im pp;
      pR=map-poly Re pp;
      g = gcd pI pR
    in
      nat ((degree p - nat (changes-R-smods-ext g (pderiv g)) - changes-R-smods
pR pI) div 2)
    )
  else
    Code.abort (STR "proots-upper fails when p=0.") (λ-. proots-upper p))
  ⟨proof⟩

```

**lemma** *proots-upper-card-code*[code]:

```

proots-upper-card p = (if p=0 then 0 else
  (let
    pf = p div (gcd p (pderiv p));
    pp = smult (inverse (lead-coeff pf)) pf;
    pI = map-poly Im pp;
    pR = map-poly Re pp;
    g = gcd pR pI
  in
    nat ((degree pf - changes-R-smods g (pderiv g) - changes-R-smods pR
pI) div 2)
  ))
  ⟨proof⟩

```

### 3.8 Polynomial roots on a general half-plane

the number of roots of polynomial  $p$ , counted with multiplicity, on the left half plane of the vector  $b - a$ .

**definition** *proots-half* :: *complex poly* ⇒ *complex* ⇒ *complex* ⇒ *nat* **where**  
*proots-half* p a b = *proots-count* p {w. Im ((w-a) / (b-a)) > 0}

**lemma** *proots-half-empty*:

```

  assumes a=b
  shows proots-half p a b = 0
  ⟨proof⟩

```

**lemma** *proots-half-proots-upper*:  
**assumes**  $a \neq b$   $p \neq 0$   
**shows**  $\text{proots-half } p \ a \ b = \text{proots-upper } (p \text{compose } p \ [ :a, (b-a):])$   
 $\langle \text{proof} \rangle$

**lemma** *proots-half-code1*[code]:  
 $\text{proots-half } p \ a \ b = (\text{if } a \neq b \text{ then}$   
      $\text{if } p \neq 0 \text{ then } \text{proots-upper } (p \circ_p \ [ :a, b - a:])$   
      $\text{else } \text{Code.abort } (\text{STR } \text{"proots-half fails when } p=0.\text{"})$   
      $(\lambda-. \text{proots-half } p \ a \ b)$   
      $\text{else } 0)$   
 $\langle \text{proof} \rangle$

### 3.9 Polynomial roots within a circle (open ball)

— Roots counted WITH multiplicity

**definition** *proots-ball*:: $\text{complex poly} \Rightarrow \text{complex} \Rightarrow \text{real} \Rightarrow \text{nat}$  **where**  
 $\text{proots-ball } p \ z0 \ r = \text{proots-count } p \ (\text{ball } z0 \ r)$

— Roots counted WITHOUT multiplicity

**definition** *proots-ball-card* :: $\text{complex poly} \Rightarrow \text{complex} \Rightarrow \text{real} \Rightarrow \text{nat}$  **where**  
 $\text{proots-ball-card } p \ z0 \ r = \text{card } (\text{proots-within } p \ (\text{ball } z0 \ r))$

**lemma** *proots-ball-code1*[code]:  
 $\text{proots-ball } p \ z0 \ r = (\text{if } r \leq 0 \text{ then}$   
      $0$   
      $\text{else if } p \neq 0 \text{ then}$   
      $\text{proots-upper } (fcompose \ (p \circ_p \ [ :z0, \text{of-real } r:]) \ [ :i, -1:] \ [ :i, 1:])$   
      $\text{else}$   
      $\text{Code.abort } (\text{STR } \text{"proots-ball fails when } p=0.\text{"})$   
      $(\lambda-. \text{proots-ball } p \ z0 \ r)$   
      $)$   
 $\langle \text{proof} \rangle$

**lemma** *proots-ball-card-code1*[code]:  
 $\text{proots-ball-card } p \ z0 \ r =$   
      $(\text{if } r \leq 0 \vee p = 0 \text{ then}$   
      $0$   
      $\text{else}$   
      $\text{proots-upper-card } (fcompose \ (p \circ_p \ [ :z0, \text{of-real } r:]) \ [ :i, -1:] \ [ :i, 1:])$   
      $)$   
 $\langle \text{proof} \rangle$

### 3.10 Polynomial roots on a circle (sphere)

— Roots counted WITH multiplicity

**definition** *proots-sphere*:: $\text{complex poly} \Rightarrow \text{complex} \Rightarrow \text{real} \Rightarrow \text{nat}$  **where**  
 $\text{proots-sphere } p \ z0 \ r = \text{proots-count } p \ (\text{sphere } z0 \ r)$

— Roots counted WITHOUT multiplicity

**definition** *proots-sphere-card* :: *complex poly*  $\Rightarrow$  *complex*  $\Rightarrow$  *real*  $\Rightarrow$  *nat* **where**  
*proots-sphere-card* *p z0 r* = *card* (*proots-within* *p* (*sphere* *z0* *r*))

**lemma** *proots-sphere-card-code1* [*code*]:

```

proots-sphere-card p z0 r =
  ( if r=0 then
    (if poly p z0=0 then 1 else 0)
    else if r < 0  $\vee$  p=0 then
      0
    else
      (if poly p (z0-r) =0 then 1 else 0) +
      proots-unbounded-line-card (fcompose (p  $\circ_p$  [z0, of-real r:]) [i, -1:])
  )

```

[*i*,1:]  
 0 1  
 )  
 <*proof*>

### 3.11 Polynomial roots on a closed ball

— Roots counted WITH multiplicity

**definition** *proots-cball*::*complex poly*  $\Rightarrow$  *complex*  $\Rightarrow$  *real*  $\Rightarrow$  *nat* **where**  
*proots-cball* *p z0 r* = *proots-count* *p* (*cball* *z0* *r*)

— Roots counted WITHOUT multiplicity

**definition** *proots-cball-card* ::*complex poly*  $\Rightarrow$  *complex*  $\Rightarrow$  *real*  $\Rightarrow$  *nat* **where**  
*proots-cball-card* *p z0 r* = *card* (*proots-within* *p* (*cball* *z0* *r*))

**lemma** *proots-cball-card-code1* [*code*]:

```

proots-cball-card p z0 r =
  ( if r=0 then
    (if poly p z0=0 then 1 else 0)
    else if r < 0  $\vee$  p=0 then
      0
    else
      ( let pp=fcompose (p  $\circ_p$  [z0, of-real r:]) [i, -1:] [i,1:]
        in
          (if poly p (z0-r) =0 then 1 else 0)
          + proots-unbounded-line-card pp 0 1
          + proots-upper-card pp
      )
  )

```

<*proof*>

**end**

## 4 Some examples for complex root counting

**theory** *Count-Complex-Roots-Examples*

```

imports Count-Complex-Roots
begin

value proots-rectangle [:2*i,0,i:] (Complex (-1) 0) (Complex 2 2)

value proots-rectangle [-1,-2*i,1:]
      (Complex (-1) 0) (Complex 2 2)

value proots-half [:1-i,2-i,1:]
      0 (Complex 0 1)

value proots-half [:1-i,2-i,1:] (Complex 0 1) 0

value [code] proots-ball ([:-2,1:]*[:-2,1:]*[:-3,1:]) 0 4

value [code] proots-ball-card ([:-2,1:]*[:-2,1:]*[:-3,1:]) 0 3

end

```

## References

- [1] M. Eisermann. The fundamental theorem of algebra made effective: An elementary real-algebraic proof via Sturm chains. *American Mathematical Monthly*, 119(9):715–752, 2012.
- [2] Q. I. Rahman and G. Schmeisser. *Analytic theory of polynomials*. Number 26. Oxford University Press, 2002.