

Computational p -Adics

Jeremy Sylvestre University of Alberta (Augustana Campus)

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Abstract

We develop a framework for reasoning computationally about p -adic fields via formal Laurent series with integer coefficients. We define a ring-class quotient type consisting of equivalence classes of prime-indexed sequences of formal Laurent series, in which every p -adic field is contained as a subring. Via a further quotient of the ring of p -adic integers (that is, the elements of depth at least 0) by the prime ideal (that is, the subring of elements of depth at least 1), we define a ring-class type in which every prime-order finite field is contained as a subring. Finally, some topological properties are developed using depth as a proxy for a metric, Hensel's Lemma is proved, and the compactness of local rings of integers is established.

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theory *Distinguished-Quotients*

imports *Main*

begin

1 Distinguished Algebraic Quotients

We define quotient groups and quotient rings modulo distinguished normal subgroups and ideals, respectively, leading to the construction of a field as a commutative ring with 1 modulo a maximal ideal.

class *distinguished-subset* =
fixes *distinguished-subset* :: 'a set ($\mathcal{N}$)

hide-const *distinguished-subset*

1.1 Quotient groups

1.1.1 Class definitions

Here we set up subclasses of *group-add* with a distinguished subgroup.

class *group-add-with-distinguished-subgroup* = *group-add* + *distinguished-subset* +
assumes *distset-nonempty*: $\mathcal{N} \neq \{\}$
and *distset-diff-closed*:
 $g \in \mathcal{N} \implies h \in \mathcal{N} \implies g - h \in \mathcal{N}$
begin

lemma *distset-zero-closed* : $0 \in \mathcal{N}$
 $\langle proof \rangle$

lemma *distset-uminus-closed*: $a \in \mathcal{N} \implies -a \in \mathcal{N}$
 $\langle proof \rangle$

lemma *distset-add-closed*:
 $a \in \mathcal{N} \implies b \in \mathcal{N} \implies a + b \in \mathcal{N}$
 $\langle proof \rangle$

lemma *distset-uminus-plus-closed*:
 $a \in \mathcal{N} \implies b \in \mathcal{N} \implies -a + b \in \mathcal{N}$
 $\langle proof \rangle$

lemma *distset-lcoset-closed1*:

```

    a ∈ N ⇒ -a + b ∈ N ⇒ b ∈ N
    ⟨proof⟩

lemma distset-lcoset-closed2:
    b ∈ N ⇒ -a + b ∈ N ⇒ a ∈ N
    ⟨proof⟩

lemma reflp-lcosetrel: - x + x ∈ N
    ⟨proof⟩

lemma symp-lcosetrel: - a + b ∈ N ⇒ - b + a ∈ N
    ⟨proof⟩

lemma transp-lcosetrel:
    -x + z ∈ N if -x + y ∈ N and -y + z ∈ N
    ⟨proof⟩

end

class group-add-with-distinguished-normal-subgroup = group-add-with-distinguished-subgroup
+
    assumes distset-normal: ((+) g) ‘ N = (λx. x + g) ‘ N

class ab-group-add-with-distinguished-subgroup =
    ab-group-add + group-add-with-distinguished-subgroup
begin

subclass group-add-with-distinguished-normal-subgroup
    ⟨proof⟩

end

```

1.1.2 The quotient group type

Here we define a quotient type relative to the left-coset relation, and instantiate it as a *group-add*.

```

quotient-type (overloaded) 'a coset-by-dist-sbgrp =
    'a::group-add-with-distinguished-subgroup / λg h :: 'a. -g + h ∈ (N::'a set)
    morphisms distinguished-quotient-coset-rep distinguished-quotient-coset
    ⟨proof⟩

```

```

lemmas distinguished-quotient-coset-rep-inverse =
    Quotient-abs-rep[OF Quotient-coset-by-dist-sbgrp]
and coset-by-dist-sbgrp-eq-iff =
    Quotient-rel-rep[OF Quotient-coset-by-dist-sbgrp]
and distinguished-quotient-coset-eqI =
    Quotient-rel-abs[OF Quotient-coset-by-dist-sbgrp]
and eq-to-distinguished-quotient-cosetI =
    Quotient-rel-abs2[OF Quotient-coset-by-dist-sbgrp]

```

lemma *coset-coset-by-dist-sbgrp-cases*

[*case-names coset, cases type: coset-by-dist-sbgrp*]:
 $(\bigwedge y. z = \text{distinguished-quotient-coset } y \implies P) \implies P$
 ⟨*proof*⟩

lemmas *distinguished-quotient-coset-eq-iff* = *coset-by-dist-sbgrp.abs-eq-iff*

lemma *distinguished-quotient-coset-rep-coset-equiv*:

– *distinguished-quotient-coset-rep* (*distinguished-quotient-coset* *r*) + *r* ∈ $\mathcal{N}$
 ⟨*proof*⟩

instantiation *coset-by-dist-sbgrp* ::

(*group-add-with-distinguished-normal-subgroup*) *group-add*

begin

lift-definition *zero-coset-by-dist-sbgrp* :: 'a *coset-by-dist-sbgrp*

is 0::'a ⟨*proof*⟩

lemmas *zero-distinguished-quotient* [simp] = *zero-coset-by-dist-sbgrp.abs-eq*[*symmetric*]

lemma *zero-distinguished-quotient-eq-iffs*:

(*distinguished-quotient-coset* *x* = 0) = (*x* ∈ $\mathcal{N}$)
 (*X* = 0) = (*distinguished-quotient-coset-rep* *X* ∈ $\mathcal{N}$)
 ⟨*proof*⟩

lift-definition *plus-coset-by-dist-sbgrp* ::

'a *coset-by-dist-sbgrp* ⇒ 'a *coset-by-dist-sbgrp* ⇒
 'a *coset-by-dist-sbgrp*

is λ*x y*::'a. *x* + *y*

⟨*proof*⟩

lemmas *plus-coset-by-dist-sbgrp* [simp] = *plus-coset-by-dist-sbgrp.abs-eq*

lift-definition *uminus-coset-by-dist-sbgrp* ::

'a *coset-by-dist-sbgrp* ⇒ 'a *coset-by-dist-sbgrp*

is λ*x*::'a. −*x*

⟨*proof*⟩

lemmas *uminus-coset-by-dist-sbgrp* [simp] =

uminus-coset-by-dist-sbgrp.abs-eq

lift-definition *minus-coset-by-dist-sbgrp* ::

'a *coset-by-dist-sbgrp* ⇒ 'a *coset-by-dist-sbgrp* ⇒
 'a *coset-by-dist-sbgrp*

is λ*x y*::'a. *x* − *y*

⟨*proof*⟩

lemmas *minus-coset-by-dist-sbgrp* [simp] = *minus-coset-by-dist-sbgrp.abs-eq*

instance $\langle proof \rangle$

end

instance *coset-by-dist-sbgrp* ::
 (*ab-group-add-with-distinguished-subgroup*) *ab-group-add*
 $\langle proof \rangle$

1.2 Quotient rings

1.2.1 Class definitions

Now we set up subclasses of *ring* with a distinguished ideal.

class *ring-with-distinguished-ideal* = *ring* + *ab-group-add-with-distinguished-subgroup* +
 assumes *distset-leftideal* : $a \in \mathcal{N} \implies r * a \in \mathcal{N}$
 and *distset-rightideal* : $a \in \mathcal{N} \implies a * r \in \mathcal{N}$

class *ring-1-with-distinguished-ideal* = *ring-1* + *ring-with-distinguished-ideal* +
 assumes *distset-no1* : $1 \notin \mathcal{N}$
 — For unitary rings we do not allow the quotient to be trivial to ensure that 0 and 1 will be different.

1.2.2 The quotient ring type

We just reuse the *'a coset-by-dist-sbgrp* type to define a quotient ring, and instantiate it as a *ring*.

type-synonym *'a distinguished-quotient-ring* = *'a coset-by-dist-sbgrp*
 — This is intended to be used with *ring-with-distinguished-ideal* or one of its subclasses.

instantiation *coset-by-dist-sbgrp* :: (*ring-with-distinguished-ideal*) *ring*
begin

lift-definition *times-coset-by-dist-sbgrp* ::
 'a distinguished-quotient-ring $\Rightarrow$ *'a distinguished-quotient-ring* $\Rightarrow$
 'a distinguished-quotient-ring
 is $\lambda x y :: 'a. x * y$
 $\langle proof \rangle$

lemmas *times-coset-by-dist-sbgrp* [*simp*] = *times-coset-by-dist-sbgrp.abs-eq*

instance $\langle proof \rangle$

end

instantiation *coset-by-dist-sbgrp* :: (*ring-1-with-distinguished-ideal*) *ring-1*

begin

lift-definition *one-coset-by-dist-sbgrp* :: 'a distinguished-quotient-ring
is 1::'a <proof>

lemmas *one-coset-by-dist-sbgrp* [simp] = *one-coset-by-dist-sbgrp.abs-eq*[symmetric]

instance <proof>

end

1.3 Quotients of commutative rings

For a commutative ring, the quotient modulo a prime ideal is an integral domain, and the quotient modulo a maximal ideal is a field.

1.3.1 Class definitions

class *comm-ring-with-distinguished-ideal* = *comm-ring* + *ab-group-add-with-distinguished-subgroup*
 +

assumes *distset-ideal* : $a \in \mathcal{N} \implies r * a \in \mathcal{N}$

begin

subclass *ring-with-distinguished-ideal*
 <proof>

definition *distideal-extension* :: 'a $\Rightarrow$ 'a set

where *distideal-extension* $a \equiv \{x. \exists r z. z \in \mathcal{N} \wedge x = r * a + z\}$

lemma *distideal-extension*: $\mathcal{N} \subseteq \text{distideal-extension } a$
 <proof>

lemma *distideal-extension-diff-closed*:

$r \in \text{distideal-extension } a \implies s \in \text{distideal-extension } a \implies$
 $r - s \in \text{distideal-extension } a$

<proof>

lemma *distideal-extension-ideal*:

$x \in \text{distideal-extension } a \implies r * x \in \text{distideal-extension } a$
 <proof>

end

class *comm-ring-1-with-distinguished-ideal* = *ring-1* + *comm-ring-with-distinguished-ideal*
 +

assumes *distset-no1*:

$1 \notin \mathcal{N}$

— Don't allow the quotient to be trivial.

begin

```

subclass ring-1-with-distinguished-ideal <proof>

lemma distideal-extension-by:  $a \in \text{distideal-extension } a$ 
  <proof>

end

class comm-ring-1-with-distinguished-pideal = comm-ring-1-with-distinguished-ideal
+
  assumes distset-pideal:  $r * s \in \mathcal{N} \implies r \in \mathcal{N} \vee s \in \mathcal{N}$ 

class comm-ring-1-with-distinguished-maxideal = comm-ring-1-with-distinguished-ideal
+
  assumes distset-maxideal:
    [
       $\mathcal{N} \subseteq I; 1 \notin I;$ 
       $\forall r s. r \in I \wedge s \in I \longrightarrow r - s \in I;$ 
       $\forall a r. a \in I \longrightarrow r * a \in I$ 
    ]  $\implies I = \mathcal{N}$ 
begin

lemma distset-maxideal-vs-distideal-extension:
   $1 \notin \text{distideal-extension } a \implies \text{distideal-extension } a = \mathcal{N}$ 
  <proof>

subclass comm-ring-1-with-distinguished-pideal
  <proof>

end

```

1.3.2 Instances and instantiations

```

instance coset-by-dist-sbgrp :: (comm-ring-with-distinguished-ideal) comm-ring
  <proof>

instance coset-by-dist-sbgrp :: (comm-ring-1-with-distinguished-ideal) comm-ring-1
  <proof>

instance coset-by-dist-sbgrp :: (comm-ring-1-with-distinguished-pideal) idom
  <proof>

lemma inverse-in-quotient-mod-maxideal:
   $\exists y. -(y * x) + 1 \in \mathcal{N}$  if  $x \notin \mathcal{N}$ 
  for  $x :: 'a :: \text{comm-ring-1-with-distinguished-maxideal}$ 
  <proof>

instantiation coset-by-dist-sbgrp :: (comm-ring-1-with-distinguished-maxideal) field
begin

```

lift-definition *inverse-coset-by-dist-sbgrp* ::
'a distinguished-quotient-ring $\Rightarrow$ 'a distinguished-quotient-ring
is $\lambda x::'a.$ if $x \in \mathcal{N}$ then 0 else (SOME $y::'a.$ $-(y*x) + 1 \in \mathcal{N}$)
 $\langle \text{proof} \rangle$

lemma *inverse-coset-by-dist-sbgrp-eq0* [simp]:
 $x \in \mathcal{N} \Longrightarrow \text{inverse} (\text{distinguished-quotient-coset } x) = 0$
 $\langle \text{proof} \rangle$

lemma *coset-by-dist-sbgrp-inverse-rep*:
 $x \notin \mathcal{N} \Longrightarrow -(y * x) + 1 \in \mathcal{N} \Longrightarrow$
 $\text{inverse} (\text{distinguished-quotient-coset } x) = \text{distinguished-quotient-coset } y$
 $\langle \text{proof} \rangle$

definition *divide-coset-by-dist-sbgrp-def* [simp]:
 $x \text{ div } y = (x::'a \text{ distinguished-quotient-ring}) * \text{inverse } y$

lemma *coset-by-dist-sbgrp-divide-rep*:
 $y \notin \mathcal{N} \Longrightarrow -(z * y) + 1 \in \mathcal{N} \Longrightarrow$
 $\text{distinguished-quotient-coset } x \text{ div distinguished-quotient-coset } y$
 $= \text{distinguished-quotient-coset } (x * z)$
 $\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$

end

end

theory *Polynomial-Extras*

imports *HOL-Computational-Algebra.Polynomial*

begin

2 Additional facts about polynomials

2.1 General facts

lemma *set-strip-while*: $\text{set} (\text{strip-while } P \text{ } xs) \subseteq \text{set } xs$
 $\langle \text{proof} \rangle$

lemma *pCons-induct'* [case-names 0 pCons0 pCons]:
 $P \text{ } p$
if zero : $P \text{ } 0$

and $pCons0 : \bigwedge a. a \neq 0 \implies P (pCons a 0)$
and $pCons-n0 : \bigwedge a p. p \neq 0 \implies P p \implies P (pCons a p)$
 $\langle proof \rangle$

lemma $pCons-decomp$: $pCons a p = pCons a 0 + pCons 0 p$
 $\langle proof \rangle$

lemma $coeffs-smult-eq-smult-coeffs$:
 $coeffs (smult x p) = map ((*) x) (strip-while (\lambda c. x * c = 0) (coeffs p))$
 $\langle proof \rangle$

lemma $coeffs-smult-eq-smult-coeffs-no-zero-divisors$:
 $coeffs (smult x p) = map ((*) x) (coeffs p)$ **if** $x \neq 0$
for $x :: 'a :: \{comm-semiring-0, semiring-no-zero-divisors\}$ **and** $p :: 'a poly$
 $\langle proof \rangle$

2.2 Derivatives

The type sort on $pderiv$ is overly restrictive for algebraic purposes, so here is the relevant material with the type sort relaxed.

function $polyderiv :: ('a :: comm-monoid-add) poly \Rightarrow 'a poly$
where $polyderiv (pCons a p) = (if p = 0 then 0 else p + pCons 0 (polyderiv p))$
 $\langle proof \rangle$

termination $polyderiv$
 $\langle proof \rangle$

declare $polyderiv.simps[simp del]$

lemma $polyderiv-0 [simp]$: $polyderiv 0 = 0$
 $\langle proof \rangle$

lemma $polyderiv-pCons$: $polyderiv (pCons a p) = p + pCons 0 (polyderiv p)$
 $\langle proof \rangle$

lemma $polyderiv-1 [simp]$: $polyderiv 1 = 0$
 $\langle proof \rangle$

2.3 Additive polynomials

The Binomial Theorem implies that to every polynomial can be associated a finite list of polynomials that describe how the original polynomial evaluates on binomial arguments.

function $additive-poly-poly :: 'a :: comm-semiring-1 poly \Rightarrow 'a poly poly$
where
 $additive-poly-poly (pCons a p) =$
 $(if p = 0$
 $then [[:a:]]$

$$\text{else smult } (pCons \ 0 \ 1) \ (additive\text{-}poly\text{-}poly \ p) + pCons \ [:a:] \ (additive\text{-}poly\text{-}poly \ p)$$

$$\langle proof \rangle$$

termination *additive-poly-poly*
 $\langle proof \rangle$

declare *additive-poly-poly.simps*[simp del]

lemma *additive-poly-poly-0* [simp]: *additive-poly-poly 0 = 0*
 $\langle proof \rangle$

lemma *additive-poly-poly-pCons0*:
 $a \neq 0 \implies additive\text{-}poly\text{-}poly \ [:a:] = [[:a:]]$
 $\langle proof \rangle$

lemma *additive-poly-poly-pCons*:
 $p \neq 0 \implies$
 $additive\text{-}poly\text{-}poly \ (pCons \ a \ p) =$
 $smult \ (pCons \ 0 \ 1) \ (additive\text{-}poly\text{-}poly \ p) + pCons \ [:a:] \ (additive\text{-}poly\text{-}poly \ p)$
 $\langle proof \rangle$

lemma *additive-poly-poly*: $poly \ p \ (x + y) = poly \ (poly \ (additive\text{-}poly\text{-}poly \ p) \ [:x:])$
 y
 $\langle proof \rangle$

lemma *additive-poly-poly-coeff0*: $coeff \ (additive\text{-}poly\text{-}poly \ p) \ 0 = p$
 $\langle proof \rangle$

lemma *additive-poly-poly-coeff1*: $coeff \ (additive\text{-}poly\text{-}poly \ p) \ 1 = polyderiv \ p$
 $\langle proof \rangle$

lemma *additive-poly-poly-eq0-iff*: $additive\text{-}poly\text{-}poly \ p = 0 \longleftrightarrow p = 0$
 $\langle proof \rangle$

lemma *additive-poly-poly-degree*: $degree \ (additive\text{-}poly\text{-}poly \ p) = degree \ p$
 $\langle proof \rangle$

lemma *poly-additive-poly-poly-pCons*:
 $poly \ (additive\text{-}poly\text{-}poly \ (pCons \ a \ p)) \ [:x:] =$
 $pCons \ a \ (poly \ (additive\text{-}poly\text{-}poly \ p) \ [:x:]) + smult \ x \ (poly \ (additive\text{-}poly\text{-}poly \ p) \ [:x:])$
if $p \neq 0$
 $\langle proof \rangle$

lemma *poly-additive-poly-poly-eq0-iff*:
 $poly \ (additive\text{-}poly\text{-}poly \ p) \ [:x:] = 0 \longleftrightarrow p = 0$
 $\langle proof \rangle$

lemma *degree-poly-additive-poly-poly*:
 $\text{degree } (\text{poly } (\text{additive-poly-poly } p) [:x:]) = \text{degree } p$
 $\langle \text{proof} \rangle$

lemma *coeff-poly-additive-poly-poly*:
 $\text{coeff } (\text{poly } (\text{additive-poly-poly } p) [:x:]) \ n = \text{poly } (\text{coeff } (\text{additive-poly-poly } p) \ n) \ x$
 $\langle \text{proof} \rangle$

lemma *additive-poly-decomp*:
 $\text{poly } p \ (x + y) = (\sum_{i \leq \text{degree } p} \text{poly } (\text{coeff } (\text{additive-poly-poly } p) \ i) \ x) * y \wedge i$
 $\langle \text{proof} \rangle$

end

theory *Parametric-Equiv-Depth*

imports
HOL.Rat
HOL-Computational-Algebra.Fraction-Field
HOL-Computational-Algebra.Primes
HOL-Library.Function-Algebras
HOL-Library.Set-Algebras
HOL.Topological-Spaces
Polynomial-Extras

begin

3 Common structures for adelic types

3.1 Preliminaries

lemma *ex-least-nat-le'*:
fixes $P :: \text{nat} \Rightarrow \text{bool}$
assumes $P \ n$
shows $\exists k \leq n. (\forall i < k. \neg P \ i) \wedge P \ k$
 $\langle \text{proof} \rangle$

lemma *ex-ex1-least-nat-le*:
fixes $P :: \text{nat} \Rightarrow \text{bool}$
assumes $\exists n. P \ n$
shows $\exists ! k. P \ k \wedge (\forall i. P \ i \longrightarrow k \leq i)$
 $\langle \text{proof} \rangle$

lemma *least-le-least*:
 $\text{Least } Q \leq \text{Least } P$

```

if  $P : \exists !x. P\ x \wedge (\forall y. P\ y \longrightarrow x \leq y)$ 
and  $Q : \exists !x. Q\ x \wedge (\forall y. Q\ y \longrightarrow x \leq y)$ 
and  $PQ: \forall x. P\ x \longrightarrow Q\ x$ 
for  $P\ Q :: 'a::order \Rightarrow bool$ 
 $\langle proof \rangle$ 

```

```

lemma linorder-wlog' [case-names le sym]:
   $\llbracket$ 
     $\bigwedge a\ b. f\ a \leq f\ b \Longrightarrow P\ a\ b;$ 
     $\bigwedge a\ b. P\ b\ a \Longrightarrow P\ a\ b$ 
   $\rrbracket \Longrightarrow P\ a\ b$ 
for  $f :: 'a \Rightarrow 'b::linorder$ 
 $\langle proof \rangle$ 

```

```

lemma log-powi-cancel [simp]:
   $a > 0 \Longrightarrow a \neq 1 \Longrightarrow \log\ a\ (a\ powi\ b) = b$ 
 $\langle proof \rangle$ 

```

3.2 Existence of primes

```

class factorial-comm-ring = factorial-semiring + comm-ring
begin subclass comm-ring-1  $\langle proof \rangle$  end

```

```

class factorial-idom = factorial-semiring + idom
begin subclass factorial-comm-ring  $\langle proof \rangle$  end

```

```

class nontrivial-factorial-semiring = factorial-semiring +
  assumes nontrivial-elem-exists:  $\exists x. x \neq 0 \wedge \text{normalize}\ x \neq 1$ 
begin

```

```

lemma primeE:
  obtains  $p$  where prime  $p$ 
 $\langle proof \rangle$ 

```

```

lemma primes-exist:  $\exists p. \text{prime}\ p$ 
 $\langle proof \rangle$ 

```

```

end

```

```

class nontrivial-factorial-comm-ring = nontrivial-factorial-semiring + comm-ring
begin
  subclass factorial-comm-ring  $\langle proof \rangle$ 
end

```

```

class nontrivial-factorial-idom = nontrivial-factorial-semiring + idom
begin
  subclass nontrivial-factorial-comm-ring  $\langle proof \rangle$ 
  subclass factorial-idom  $\langle proof \rangle$ 
end

```

```

instance int :: nontrivial-factorial-idom
  ⟨proof⟩

typedef (overloaded) 'a prime =
  {p::'a::nontrivial-factorial-semiring. prime p}
  ⟨proof⟩

lemma Rep-prime-n0: Rep-prime p ≠ 0
  ⟨proof⟩

lemma Rep-prime-not-unit: ¬ is-unit (Rep-prime p)
  ⟨proof⟩

abbreviation pdvd ::
  'a::nontrivial-factorial-semiring prime ⇒ 'a ⇒ bool (infix pdvd 50)
  where p pdvd a ≡ (Rep-prime p) dvd a
abbreviation pmultiplicity p ≡ multiplicity (Rep-prime p)

lemma multiplicity-distinct-primes:
  p ≠ q ⇒ pmultiplicity p (Rep-prime q) = 0
  ⟨proof⟩

```

3.3 Generic algebraic equality

```

context
  fixes R
  assumes sympR : symp R
  and transpR: transp R
begin

lemma transp-left: R x z ⇒ R y z ⇒ R x y
  ⟨proof⟩

lemma transp-right: R x y ⇒ R x z ⇒ R y z
  ⟨proof⟩

lemma transp-iffs:
  assumes R x y shows R x z ⟷ R y z and R z x ⟷ R z y
  ⟨proof⟩

lemma transp-cong: R w z if R x w and R y z and R x y
  ⟨proof⟩

lemma transp-cong-sym: R x y if R x w and R y z and R w z
  ⟨proof⟩

end

```

lemma *equivp-imp-symp*: $\text{equivp } R \implies \text{symp } R$
 <proof>

locale *generic-alg-equality* =
 fixes *alg-eq* :: '*a* $\Rightarrow$ '*a* $\Rightarrow$ bool
 assumes *equivp*: *equivp alg-eq*
begin

lemmas *refl* = *equivp-reflp* [*OF equivp*]
and *sym* [*sym*] = *equivp-symp* [*OF equivp*]
and *trans* [*trans*] = *equivp-transp* [*OF equivp*]
and *trans-left* = *transp-left* [*OF equivp-imp-symp equivp-imp-transp, OF equivp equivp*]
and *trans-right* = *transp-right* [*OF equivp-imp-symp equivp-imp-transp, OF equivp equivp*]
and *trans-iffs* = *transp-iffs* [*OF equivp-imp-symp equivp-imp-transp, OF equivp equivp*]
and *cong* = *transp-cong* [*OF equivp-imp-symp equivp-imp-transp, OF equivp equivp*]
and *cong-sym* = *transp-cong-sym* [*OF equivp-imp-symp equivp-imp-transp, OF equivp equivp*]

lemma *sym-iff*: $\text{alg-eq } x \ y = \text{alg-eq } y \ x$
 <proof>

end

locale *ab-group-alg-equality* = *generic-alg-equality alg-eq*
for *alg-eq* :: '*a*::*ab-group-add* $\Rightarrow$ '*a* $\Rightarrow$ bool
 +
 assumes *conv-0*: $\text{alg-eq } x \ y \longleftrightarrow \text{alg-eq } (x - y) \ 0$
begin

lemmas *trans0-iff* = *trans-iffs*(1)[*of - - 0*]

lemma *add-equiv0-left*: $\text{alg-eq } (x + y) \ y \ \text{if } \text{alg-eq } x \ 0$
 <proof>

lemma *add-equiv0-right*: $\text{alg-eq } (x + y) \ x \ \text{if } \text{alg-eq } y \ 0$
 <proof>

lemma *add*:
 $\text{alg-eq } (x1 + x2) \ (y1 + y2) \ \text{if } \text{alg-eq } x1 \ y1 \ \text{and } \text{alg-eq } x2 \ y2$
 <proof>

lemma *add-left*: $\text{alg-eq } (x + y) \ (x + y') \ \text{if } \text{alg-eq } y \ y'$
 <proof>

lemma *add-right*: $\text{alg-eq } (x + y) (x' + y) \text{ if } \text{alg-eq } x x'$
 $\langle \text{proof} \rangle$

lemma *add-0-left*: $\text{alg-eq } (x + y) y \text{ if } \text{alg-eq } x 0$
 $\langle \text{proof} \rangle$

lemma *add-0-right*: $\text{alg-eq } (x + y) x \text{ if } \text{alg-eq } y 0$
 $\langle \text{proof} \rangle$

lemma *sum-zeros*:
 $\text{finite } A \implies \forall a \in A. \text{alg-eq } (f a) 0 \implies \text{alg-eq } (\text{sum } f A) 0$
 $\langle \text{proof} \rangle$

lemma *uminus*: $\text{alg-eq } x y \implies \text{alg-eq } (-x) (-y)$
 $\langle \text{proof} \rangle$

lemma *uminus-iff*:
 $\text{alg-eq } x y \longleftrightarrow \text{alg-eq } (-x) (-y)$
 $\langle \text{proof} \rangle$

lemma *add-0*: $\text{alg-eq } (x + y) 0 \text{ if } \text{alg-eq } x 0 \text{ and } \text{alg-eq } y 0$
 $\langle \text{proof} \rangle$

lemma *minus*:
 $\text{alg-eq } x1 y1 \implies \text{alg-eq } x2 y2 \implies$
 $\text{alg-eq } (x1 - x2) (y1 - y2)$
 $\langle \text{proof} \rangle$

lemma *minus-0*: $\text{alg-eq } (x - y) 0 \text{ if } \text{alg-eq } x 0 \text{ and } \text{alg-eq } y 0$
 $\langle \text{proof} \rangle$

lemma *minus-left*: $\text{alg-eq } (x - y) (x - y') \text{ if } \text{alg-eq } y y'$
 $\langle \text{proof} \rangle$

lemma *minus-right*: $\text{alg-eq } (x - y) (x' - y) \text{ if } \text{alg-eq } x x'$
 $\langle \text{proof} \rangle$

lemma *minus-0-left*: $\text{alg-eq } (x - y) (-y) \text{ if } \text{alg-eq } x 0$
 $\langle \text{proof} \rangle$

lemma *minus-0-right*: $\text{alg-eq } (x - y) x \text{ if } \text{alg-eq } y 0$
 $\langle \text{proof} \rangle$

end

locale *ring-alg-equality* = *ab-group-alg-equality* *alg-eq*
for *alg-eq* :: 'a::comm-ring $\Rightarrow$ 'a $\Rightarrow$ bool

```

+
  assumes mult-0-right: alg-eq y 0  $\implies$  alg-eq (x * y) 0
begin

lemma mult-0-left: alg-eq x 0  $\implies$  alg-eq (x * y) 0
  <proof>

lemma mult:
  alg-eq x1 y1  $\implies$  alg-eq x2 y2  $\implies$  alg-eq (x1 * x2) (y1 * y2)
  <proof>

lemma mult-left: alg-eq y y'  $\implies$  alg-eq (x * y) (x * y')
  <proof>

lemma mult-right: alg-eq x x'  $\implies$  alg-eq (x * y) (x' * y)
  <proof>

end

locale ring-1-alg-equality = ring-alg-equality alg-eq
  for alg-eq :: 'a::comm-ring-1  $\Rightarrow$  'a  $\Rightarrow$  bool
begin

lemma mult-one-left: alg-eq x 1  $\implies$  alg-eq (x * y) y
  <proof>

lemma mult-one-right: alg-eq y 1  $\implies$  alg-eq (x * y) x
  <proof>

lemma prod-ones:
  finite A  $\implies$   $\forall a \in A. \text{alg-eq } (f \ a) \ 1 \implies$ 
    alg-eq (prod f A) 1
  <proof>

lemma prod-with-zero:
  alg-eq (prod f (insert a A)) 0 if finite A and alg-eq (f a) 0
  <proof>

lemma pow: alg-eq (x ^ n) (y ^ n) if alg-eq x y
  <proof>

lemma pow-equiv0: alg-eq (x ^ n) 0 if n > 0 and alg-eq x 0
  <proof>

end

```

3.4 Indexed equality

3.4.1 General structure

```

locale p-equality =
  fixes p-equal :: 'a ⇒ 'b::comm-ring ⇒ 'b ⇒ bool
  assumes p-alg-equality: ring-alg-equality (p-equal p)
begin

lemmas p-group-alg-equality = ring-alg-equality.axioms(1)[OF p-alg-equality]
lemmas p-generic-alg-equality = ab-group-alg-equality.axioms(1)[OF p-group-alg-equality]

lemmas equivp      = generic-alg-equality.equivp      [OF p-generic-alg-equality]
and refl[simp]    = generic-alg-equality.refl        [OF p-generic-alg-equality]
and sym [sym]     = generic-alg-equality.sym         [OF p-generic-alg-equality]
and trans [trans] = generic-alg-equality.trans       [OF p-generic-alg-equality]
and sym-iff       = generic-alg-equality.sym-iff     [OF p-generic-alg-equality]
and trans-left    = generic-alg-equality.trans-left  [OF p-generic-alg-equality]
and trans-right   = generic-alg-equality.trans-right [OF p-generic-alg-equality]
and trans-iffs    = generic-alg-equality.trans-iffs  [OF p-generic-alg-equality]
and cong          = generic-alg-equality.cong        [OF p-generic-alg-equality]
and cong-sym      = generic-alg-equality.cong-sym    [OF p-generic-alg-equality]

lemmas conv-0      = ab-group-alg-equality.conv-0    [OF p-group-alg-equality]
and trans0-iff     = ab-group-alg-equality.trans0-iff [OF p-group-alg-equality]
and add-equiv0-left = ab-group-alg-equality.add-equiv0-left [OF p-group-alg-equality]
and add-equiv0-right = ab-group-alg-equality.add-equiv0-right [OF p-group-alg-equality]
and add            = ab-group-alg-equality.add        [OF p-group-alg-equality]
and add-left       = ab-group-alg-equality.add-left   [OF p-group-alg-equality]
and add-right      = ab-group-alg-equality.add-right  [OF p-group-alg-equality]
and add-0-left     = ab-group-alg-equality.add-0-left [OF p-group-alg-equality]
and add-0-right    = ab-group-alg-equality.add-0-right [OF p-group-alg-equality]
and sum-zeros      = ab-group-alg-equality.sum-zeros  [OF p-group-alg-equality]
and uminus         = ab-group-alg-equality.uminus    [OF p-group-alg-equality]
and uminus-iff     = ab-group-alg-equality.uminus-iff [OF p-group-alg-equality]
and add-0          = ab-group-alg-equality.add-0     [OF p-group-alg-equality]
and minus          = ab-group-alg-equality.minus      [OF p-group-alg-equality]
and minus-0       = ab-group-alg-equality.minus-0    [OF p-group-alg-equality]
and minus-left     = ab-group-alg-equality.minus-left [OF p-group-alg-equality]
and minus-right    = ab-group-alg-equality.minus-right [OF p-group-alg-equality]
and minus-0-left   = ab-group-alg-equality.minus-0-left [OF p-group-alg-equality]
and minus-0-right  = ab-group-alg-equality.minus-0-right [OF p-group-alg-equality]

lemmas mult-0-left = ring-alg-equality.mult-0-left [OF p-alg-equality]
and mult-0-right = ring-alg-equality.mult-0-right [OF p-alg-equality]
and mult         = ring-alg-equality.mult         [OF p-alg-equality]
and mult-left    = ring-alg-equality.mult-left    [OF p-alg-equality]
and mult-right   = ring-alg-equality.mult-right   [OF p-alg-equality]

definition globally-p-equal :: 'b ⇒ 'b ⇒ bool

```

where *globally-p-equal-def[simp]*: $\text{globally-p-equal } x \ y = (\forall p::'a. \text{p-equal } p \ x \ y)$

lemma *globally-p-equalI[intro]*: $\text{globally-p-equal } x \ y$ **if** $\bigwedge p::'a. \text{p-equal } p \ x \ y$
 ⟨proof⟩

lemma *globally-p-equalD*: $\text{globally-p-equal } x \ y \implies \text{p-equal } p \ x \ y$
 ⟨proof⟩

lemma *globally-p-nequalE*:
 assumes $\neg \text{globally-p-equal } x \ y$
 obtains p **where** $\neg \text{p-equal } p \ x \ y$
 ⟨proof⟩

sublocale *globally-p-equality*: *ring-alg-equality* *globally-p-equal*
 ⟨proof⟩

lemmas *globally-p-equal-equivp* = *globally-p-equality-equivp*
and *globally-p-equal-refl[simp]* = *globally-p-equality-refl*
and *globally-p-equal-sym* = *globally-p-equality-sym*
and *globally-p-equal-trans* = *globally-p-equality-trans*
and *globally-p-equal-trans-left* = *globally-p-equality-trans-left*
and *globally-p-equal-trans-right* = *globally-p-equality-trans-right*
and *globally-p-equal-trans-iffs* = *globally-p-equality-trans-iffs*
and *globally-p-equal-cong* = *globally-p-equality-cong*
and *globally-p-equal-conv-0* = *globally-p-equality-conv-0*
and *globally-p-equal-trans0-iff* = *globally-p-equality-trans0-iff*
and *globally-p-equal-add-equiv0-left* = *globally-p-equality-add-equiv0-left*
and *globally-p-equal-add-equiv0-right* = *globally-p-equality-add-equiv0-right*
and *globally-p-equal-add* = *globally-p-equality-add*
and *globally-p-equal-uminus* = *globally-p-equality-uminus*
and *globally-p-equal-uminus-iff* = *globally-p-equality-uminus-iff*
and *globally-p-equal-add-0* = *globally-p-equality-add-0*
and *globally-p-equal-minus* = *globally-p-equality-minus*
and *globally-p-equal-mult-0-left* = *globally-p-equality-mult-0-left*
and *globally-p-equal-mult-0-right* = *globally-p-equality-mult-0-right*
and *globally-p-equal-mult* = *globally-p-equality-mult*
and *globally-p-equal-mult-left* = *globally-p-equality-mult-left*
and *globally-p-equal-mult-right* = *globally-p-equality-mult-right*

lemma *globally-p-equal-transfer-equals-rsp*:
 $\text{p-equal } p \ x1 \ y1 \longleftrightarrow \text{p-equal } p \ x2 \ y2$
if $\text{globally-p-equal } x1 \ x2$ **and** $\text{globally-p-equal } y1 \ y2$
 ⟨proof⟩

end

locale *p-equality-1* = *p-equality* *p-equal*
for $\text{p-equal} :: 'a \Rightarrow 'b::\text{comm-ring-1} \Rightarrow 'b \Rightarrow \text{bool}$

begin

lemma *p-alg-equality-1*: *ring-1-alg-equality* (*p-equal* *p*)
 ⟨*proof*⟩

lemmas *mult-one-left* = *ring-1-alg-equality.mult-one-left* [*OF p-alg-equality-1*]
and *mult-one-right* = *ring-1-alg-equality.mult-one-right* [*OF p-alg-equality-1*]
and *prod-ones* = *ring-1-alg-equality.prod-ones* [*OF p-alg-equality-1*]
and *pow* = *ring-1-alg-equality.pow* [*OF p-alg-equality-1*]
and *pow-equiv0* = *ring-1-alg-equality.pow-equiv0* [*OF p-alg-equality-1*]

end

locale *p-equality-no-zero-divisors* = *p-equality*

+

assumes *no-zero-divisors*:

$\neg p\text{-equal } p \ x \ 0 \implies \neg p\text{-equal } p \ y \ 0 \implies$
 $\neg p\text{-equal } p \ (x * y) \ 0$

begin

lemma *mult-equiv-0-iff*:

$p\text{-equal } p \ (x * y) \ 0 \longleftrightarrow$
 $p\text{-equal } p \ x \ 0 \vee p\text{-equal } p \ y \ 0$
 ⟨*proof*⟩

end

locale *p-equality-no-zero-divisors-1* = *p-equality-no-zero-divisors* *p-equal*

for *p-equal* ::

$'a \Rightarrow 'b :: \text{comm-ring-1} \Rightarrow 'b \Rightarrow \text{bool}$

+

assumes *nontrivial*: $\exists x. \neg p\text{-equal } p \ x \ 0$

begin

sublocale *p-equality-1* ⟨*proof*⟩

lemma *one-p-nequal-zero*: $\neg p\text{-equal } p \ (1::'b) \ (0::'b)$
 ⟨*proof*⟩

lemma *zero-p-nequal-one*: $\neg p\text{-equal } p \ (0::'b) \ (1::'b)$
 ⟨*proof*⟩

lemma *neg-one-p-nequal-zero*: $\neg p\text{-equal } p \ (-1::'b) \ (0::'b)$
 ⟨*proof*⟩

lemma *pow-equiv-0-iff*: $p\text{-equal } p \ (x \wedge n) \ 0 \longleftrightarrow p\text{-equal } p \ x \ 0 \wedge n \neq 0$
 ⟨*proof*⟩

lemma *pow-equiv-base*: $p\text{-equal } p (x \wedge n) (y \wedge n)$ **if** $p\text{-equal } p x y$
 ⟨*proof*⟩

end

locale *p-equality-div-inv* = *p-equality-no-zero-divisors-1* *p-equal*
for *p-equal* ::

'*a* $\Rightarrow$ '*b*::{*comm-ring-1*, *inverse*, *divide-trivial*} $\Rightarrow$
 '*b* $\Rightarrow$ *bool*

+

assumes *divide-inverse* : $\bigwedge x y :: 'b. x / y = x * \text{inverse } y$
and *inverse-inverse* : $\bigwedge x :: 'b. \text{inverse } (\text{inverse } x) = x$
and *inverse-mult* : $\bigwedge x y :: 'b. \text{inverse } (x * y) = \text{inverse } x * \text{inverse } y$
and *inverse-equiv0-iff*: $p\text{-equal } p (\text{inverse } x) 0 \longleftrightarrow p\text{-equal } p x 0$
and *divide-self-equiv* : $\neg p\text{-equal } p x 0 \Longrightarrow p\text{-equal } p (x / x) 1$

begin

Global algebra facts

lemma *inverse-eq-divide*: $\text{inverse } x = 1 / x$ **for** $x :: 'b$
 ⟨*proof*⟩

lemma *inverse-div-inverse*: $\text{inverse } x / \text{inverse } y = y / x$ **for** $x y :: 'b$
 ⟨*proof*⟩

lemma *inverse-0[simp]*: $\text{inverse } (0 :: 'b) = 0$
 ⟨*proof*⟩

lemma *inverse-1[simp]*: $\text{inverse } (1 :: 'b) = 1$
 ⟨*proof*⟩

lemma *inverse-pow*: $\text{inverse } (x \wedge n) = \text{inverse } x \wedge n$ **for** $x :: 'b$
 ⟨*proof*⟩

lemma *power-int-minus*: $\text{power-int } x (-n) = \text{inverse } (\text{power-int } x n)$ **for** $x :: 'b$
 ⟨*proof*⟩

lemma *power-int-minus-divide*: $\text{power-int } x (-n) = 1 / (\text{power-int } x n)$ **for** $x :: 'b$
 ⟨*proof*⟩

lemma *power-int-0-left-if*: $\text{power-int } (0 :: 'b) m = (\text{if } m = 0 \text{ then } 1 \text{ else } 0)$
 ⟨*proof*⟩

lemma *power-int-0-left [simp]*: $m \neq 0 \Longrightarrow \text{power-int } (0 :: 'b) m = 0$
 ⟨*proof*⟩

lemma *power-int-1-left [simp]*: $\text{power-int } (1 :: 'b) n = 1$
 ⟨*proof*⟩

lemma *inverse-power-int*: $\text{inverse } (\text{power-int } x \ n) = \text{power-int } x \ (-n)$ **for** $x :: 'b$
 $\langle \text{proof} \rangle$

lemma *power-int-inverse*:
 $\text{power-int } (\text{inverse } x) \ n = \text{inverse } (\text{power-int } x \ n)$ **for** $x :: 'b$
 $\langle \text{proof} \rangle$

lemma *power-int-mult-distrib*:
 $\text{power-int } (x * y) \ m = \text{power-int } x \ m * \text{power-int } y \ m$ **for** $x \ y :: 'b$
 $\langle \text{proof} \rangle$

lemma *inverse-divide*: $\text{inverse } (a / b) = b / a$ **for** $a \ b :: 'b$
 $\langle \text{proof} \rangle$

lemma *minus-divide-left*: $-(a / b) = (-a) / b$ **for** $a \ b :: 'b$
 $\langle \text{proof} \rangle$

lemma *times-divide-times-eq*: $(a / b) * (c / d) = (a * c) / (b * d)$ **for** $a \ b :: 'b$
 $\langle \text{proof} \rangle$

lemma *divide-pow*: $(x \wedge n) / (y \wedge n) = (x / y) \wedge n$ **for** $x \ y :: 'b$
 $\langle \text{proof} \rangle$

lemma *power-int-divide-distrib*:
 $\text{power-int } (x / y) \ m = \text{power-int } x \ m / \text{power-int } y \ m$ **for** $x \ y :: 'b$
 $\langle \text{proof} \rangle$

lemma *power-int-one-over*: $\text{power-int } (1 / x) \ n = 1 / \text{power-int } x \ n$ **for** $x :: 'b$
 $\langle \text{proof} \rangle$

lemma *add-divide-distrib*: $(a + b) / c = (a / c) + (b / c)$ **for** $a \ b \ c :: 'b$
 $\langle \text{proof} \rangle$

lemma *diff-divide-distrib*: $(a - b) / c = (a / c) - (b / c)$ **for** $a \ b \ c :: 'b$
 $\langle \text{proof} \rangle$

lemma *times-divide-eq-right*: $a * (b / c) = (a * b) / c$ **for** $a \ b \ c :: 'b$
 $\langle \text{proof} \rangle$

lemma *times-divide-eq-left*: $(b / c) * a = (b * a) / c$ **for** $a \ b \ c :: 'b$
 $\langle \text{proof} \rangle$

lemma *divide-divide-eq-left*: $(a / b) / c = a / (b * c)$ **for** $a \ b \ c :: 'b$
 $\langle \text{proof} \rangle$

lemma *swap-numerator*: $a * (b / c) = (a / c) * b$ **for** $a \ b \ c :: 'b$
 $\langle \text{proof} \rangle$

lemma *divide-divide-times-eq*: $(a / b) / (c / d) = (a * d) / (b * c)$ **for** $a\ b\ c\ d :: 'b$

$\langle proof \rangle$

Algebra facts by place

lemma *right-inverse-equiv*: $p\text{-equal}\ p\ (x * inverse\ x)\ 1$ **if** $\neg p\text{-equal}\ p\ x\ 0$

$\langle proof \rangle$

lemma *left-inverse-equiv*: $p\text{-equal}\ p\ (inverse\ x * x)\ 1$ **if** $\neg p\text{-equal}\ p\ x\ 0$

$\langle proof \rangle$

lemma *inverse-p-unique*:

$p\text{-equal}\ p\ (inverse\ x)\ y$ **if** $p\text{-equal}\ p\ (x * y)\ 1$

$\langle proof \rangle$

lemma *inverse-equiv-imp-equiv*:

$p\text{-equal}\ p\ (inverse\ a)\ (inverse\ b) \implies p\text{-equal}\ p\ a\ b$

$\langle proof \rangle$

lemma *inverse-pcong*: $p\text{-equal}\ p\ (inverse\ x)\ (inverse\ y)$ **if** $p\text{-equal}\ p\ x\ y$

$\langle proof \rangle$

lemma *inverse-equiv-iff-equiv*:

$p\text{-equal}\ p\ (inverse\ a)\ (inverse\ b) \longleftrightarrow p\text{-equal}\ p\ a\ b$

$\langle proof \rangle$

lemma *power-int-equiv-0-iff*:

$p\text{-equal}\ p\ (power\text{-int}\ x\ n)\ 0 \longleftrightarrow p\text{-equal}\ p\ x\ 0 \wedge n \neq 0$

$\langle proof \rangle$

lemma *power-diff-conv-inverse-equiv*:

$m \leq n \implies p\text{-equal}\ p\ (x \wedge (n - m))\ (x \wedge n * inverse\ x \wedge m)$

if $\neg p\text{-equal}\ p\ x\ 0$

$\langle proof \rangle$

lemma *power-int-add-1-equiv*:

$p\text{-equal}\ p\ (power\text{-int}\ x\ (m + 1))\ (power\text{-int}\ x\ m * x)$

if $\neg p\text{-equal}\ p\ x\ 0 \vee m \neq -1$

$\langle proof \rangle$

lemma *power-int-add-equiv*:

$\neg p\text{-equal}\ p\ x\ 0 \vee m + n \neq 0 \implies$

$p\text{-equal}\ p\ (power\text{-int}\ x\ (m + n))\ (power\text{-int}\ x\ m * power\text{-int}\ x\ n)$

$\langle proof \rangle$

lemma *power-int-diff-equiv*:

$p\text{-equal}\ p\ (power\text{-int}\ x\ (m - n))\ (power\text{-int}\ x\ m / power\text{-int}\ x\ n)$

if $\neg p\text{-equal}\ p\ x\ 0 \vee m \neq n$

$\langle proof \rangle$

lemma *power-int-minus-mult-equiv*:

$p\text{-equal } p \text{ (power-int } x \text{ (} n - 1 \text{) } * x \text{) (power-int } x \text{ } n \text{)}$
if $\neg p\text{-equal } p \text{ } x \text{ } 0 \vee n \neq 0$
 $\langle \text{proof} \rangle$

lemma *power-int-not-equiv-zero*:

$\neg p\text{-equal } p \text{ } x \text{ } 0 \vee n = 0 \implies \neg p\text{-equal } p \text{ (power-int } x \text{ } n \text{) } 0$
 $\langle \text{proof} \rangle$

lemma *power-int-equiv-base*: $p\text{-equal } p \text{ (power-int } x \text{ } n \text{) (power-int } y \text{ } n \text{) if } p\text{-equal } p \text{ } x \text{ } y$

$\langle \text{proof} \rangle$

lemma *divide-equiv-0-iff*:

$p\text{-equal } p \text{ (} x \text{ / } y \text{) } 0 \iff p\text{-equal } p \text{ } x \text{ } 0 \vee p\text{-equal } p \text{ } y \text{ } 0$
 $\langle \text{proof} \rangle$

lemma *divide-pcong*: $p\text{-equal } p \text{ (} w \text{ / } x \text{) (} y \text{ / } z \text{) if } p\text{-equal } p \text{ } w \text{ } y \text{ and } p\text{-equal } p \text{ } x \text{ } z$

$\langle \text{proof} \rangle$

lemma *divide-left-equiv*: $p\text{-equal } p \text{ } x \text{ } y \implies p\text{-equal } p \text{ (} x \text{ / } z \text{) (} y \text{ / } z \text{)}$

$\langle \text{proof} \rangle$

lemma *divide-right-equiv*: $p\text{-equal } p \text{ } x \text{ } y \implies p\text{-equal } p \text{ (} z \text{ / } x \text{) (} z \text{ / } y \text{)}$

$\langle \text{proof} \rangle$

lemma *add-frac-equiv*:

$p\text{-equal } p \text{ ((} a \text{ / } b \text{) + (} c \text{ / } d \text{)) ((} a * d + c * b \text{) / (} b * d \text{))}$
if $\neg p\text{-equal } p \text{ } b \text{ } 0$ **and** $\neg p\text{-equal } p \text{ } d \text{ } 0$
 $\langle \text{proof} \rangle$

lemma *mult-divide-mult-cancel-right*:

$p\text{-equal } p \text{ ((} a * b \text{) / (} c * b \text{)) (} a \text{ / } c \text{) if } \neg p\text{-equal } p \text{ } b \text{ } 0$
 $\langle \text{proof} \rangle$

lemma *mult-divide-mult-cancel-right2*:

$p\text{-equal } p \text{ ((} a * c \text{) / (} c * b \text{)) (} a \text{ / } b \text{) if } \neg p\text{-equal } p \text{ } c \text{ } 0$
 $\langle \text{proof} \rangle$

lemma *mult-divide-mult-cancel-left*:

$p\text{-equal } p \text{ ((} c * a \text{) / (} c * b \text{)) (} a \text{ / } b \text{) if } \neg p\text{-equal } p \text{ } c \text{ } 0$
 $\langle \text{proof} \rangle$

lemma *mult-divide-mult-cancel-left2*:

$p\text{-equal } p \text{ ((} c * a \text{) / (} b * c \text{)) (} a \text{ / } b \text{) if } \neg p\text{-equal } p \text{ } c \text{ } 0$
 $\langle \text{proof} \rangle$

lemma *mult-divide-cancel-right*: $p\text{-equal } p \text{ ((} a * b \text{) / } b \text{) } a \text{ if } \neg p\text{-equal } p \text{ } b \text{ } 0$

$\langle \text{proof} \rangle$

lemma *mult-divide-cancel-left*: $p\text{-equal } p ((a * b) / a) b \text{ if } \neg p\text{-equal } p a 0$
 $\langle \text{proof} \rangle$

lemma *divide-equiv-equiv*: $(p\text{-equal } p (b / c) a) = (p\text{-equal } p b (a * c)) \text{ if } \neg p\text{-equal } p c 0$
 $\langle \text{proof} \rangle$

lemma *neg-divide-equiv-equiv*:
 $p\text{-equal } p (- (b / c)) a \iff p\text{-equal } p (-b) (a * c) \text{ if } \neg p\text{-equal } p c 0$
 $\langle \text{proof} \rangle$

lemma *equiv-divide-imp*:
 $p\text{-equal } p (a * c) b \implies p\text{-equal } p a (b / c) \text{ if } \neg p\text{-equal } p c 0$
 $\langle \text{proof} \rangle$

lemma *add-divide-equiv-iff*:
 $p\text{-equal } p (x + y / z) ((x * z + y) / z) \text{ if } \neg p\text{-equal } p z 0$
 $\langle \text{proof} \rangle$

lemma *divide-add-equiv-iff*: $p\text{-equal } p (x / z + y) ((x + y * z) / z) \text{ if } \neg p\text{-equal } p z 0$
 $\langle \text{proof} \rangle$

lemma *diff-divide-equiv-iff*: $p\text{-equal } p (x - y / z) ((x * z - y) / z) \text{ if } \neg p\text{-equal } p z 0$
 $\langle \text{proof} \rangle$

lemma *minus-divide-add-equiv-iff*:
 $p\text{-equal } p (- (x / z) + y) ((- x + y * z) / z) \text{ if } \neg p\text{-equal } p z 0$
 $\langle \text{proof} \rangle$

lemma *divide-diff-equiv-iff*: $p\text{-equal } p (x / z - y) ((x - y * z) / z) \text{ if } \neg p\text{-equal } p z 0$
 $\langle \text{proof} \rangle$

lemma *minus-divide-diff-equiv-iff*:
 $p\text{-equal } p (- (x / z) - y) ((- x - y * z) / z) \text{ if } \neg p\text{-equal } p z 0$
 $\langle \text{proof} \rangle$

lemma *divide-mult-cancel-left*: $p\text{-equal } p (a / (a * b)) (1 / b) \text{ if } \neg p\text{-equal } p a 0$
 $\langle \text{proof} \rangle$

lemma *divide-mult-cancel-right*: $p\text{-equal } p (b / (a * b)) (1 / a) \text{ if } \neg p\text{-equal } p b 0$
 $\langle \text{proof} \rangle$

lemma *diff-frac-equiv*:
 $p\text{-equal } p (x / y - w / z) ((x * z - w * y) / (y * z))$

if $\neg p\text{-equal } p \ y \ 0$ **and** $\neg p\text{-equal } p \ z \ 0$
 $\langle \text{proof} \rangle$

lemma *frac-equiv-equiv*:
 $p\text{-equal } p \ (x \ / \ y) \ (w \ / \ z) \longleftrightarrow p\text{-equal } p \ (x \ * \ z) \ (w \ * \ y)$
if $\neg p\text{-equal } p \ y \ 0$ **and** $\neg p\text{-equal } p \ z \ 0$
 $\langle \text{proof} \rangle$

lemma *inverse-equiv-1-iff*:
 $p\text{-equal } p \ (\text{inverse } x) \ 1 \longleftrightarrow p\text{-equal } p \ x \ 1$
 $\langle \text{proof} \rangle$

lemma *inverse-neg-1-equiv*: $p\text{-equal } p \ (\text{inverse } (-1)) \ (-1)$
 $\langle \text{proof} \rangle$

lemma *globally-inverse-neg-1*: $\text{globally-}p\text{-equal } (\text{inverse } (-1)) \ (-1)$
 $\langle \text{proof} \rangle$

lemma *inverse-uminus-equiv*: $p\text{-equal } p \ (\text{inverse } (-x)) \ (- \ \text{inverse } x)$
 $\langle \text{proof} \rangle$

lemma *minus-divide-right-equiv*: $p\text{-equal } p \ (a \ / \ (- \ b)) \ (- \ (a \ / \ b))$
 $\langle \text{proof} \rangle$

lemma *minus-divide-divide-equiv*: $p\text{-equal } p \ ((- \ a) \ / \ (- \ b)) \ (a \ / \ b)$
 $\langle \text{proof} \rangle$

lemma *divide-minus1-equiv*: $p\text{-equal } p \ (x \ / \ (-1)) \ (- \ x)$
 $\langle \text{proof} \rangle$

lemma *divide-cancel-right*:
 $p\text{-equal } p \ (a \ / \ c) \ (b \ / \ c) \longleftrightarrow p\text{-equal } p \ c \ 0 \ \vee \ p\text{-equal } p \ a \ b$
 $\langle \text{proof} \rangle$

lemma *divide-cancel-left*:
 $p\text{-equal } p \ (c \ / \ a) \ (c \ / \ b) \longleftrightarrow p\text{-equal } p \ c \ 0 \ \vee \ p\text{-equal } p \ a \ b$
 $\langle \text{proof} \rangle$

lemma *divide-equiv-1-iff*:
 $p\text{-equal } p \ (a \ / \ b) \ 1 \longleftrightarrow \neg p\text{-equal } p \ b \ 0 \ \wedge \ p\text{-equal } p \ a \ b$
 $\langle \text{proof} \rangle$

lemma *divide-equiv-minus-1-iff*:
 $p\text{-equal } p \ (a \ / \ b) \ (- \ 1) \longleftrightarrow \neg p\text{-equal } p \ b \ 0 \ \wedge \ p\text{-equal } p \ a \ (- \ b)$
 $\langle \text{proof} \rangle$

end

3.5 Indexed depth functions

3.5.1 General structure

```

locale p-equality-depth = p-equality +
  fixes p-depth :: 'a  $\Rightarrow$  'b::comm-ring  $\Rightarrow$  int
  assumes depth-of-0[simp]: p-depth p 0 = 0
  and    depth-equiv: p-equal p x y  $\implies$  p-depth p x = p-depth p y
  and    depth-uminus: p-depth p ( $-x$ ) = p-depth p x
  and    depth-pre-nonarch:
     $\neg$  p-equal p x 0  $\implies$  p-depth p x < p-depth p y  $\implies$ 
      p-depth p (x + y) = p-depth p x
     $\llbracket \neg$  p-equal p x 0;  $\neg$  p-equal p (x + y) 0; p-depth p x = p-depth p y  $\rrbracket$ 
       $\implies$  p-depth p (x + y)  $\geq$  p-depth p x
begin

lemma depth-equiv-0: p-equal p x 0  $\implies$  p-depth p x = 0
  <proof>

lemma depth-equiv-uminus: p-equal p y ( $-x$ )  $\implies$  p-depth p y = p-depth p x
  <proof>

lemma depth-add-equiv0-left: p-equal p x 0  $\implies$  p-depth p (x + y) = p-depth p y
  <proof>

lemma depth-add-equiv0-right: p-equal p y 0  $\implies$  p-depth p (x + y) = p-depth p x
  <proof>

lemma depth-add-equiv:
  p-depth p (x + y) = p-depth p (w + z) if p-equal p x w and p-equal p y z
  <proof>

lemma depth-diff: p-depth p (x - y) = p-depth p (y - x)
  <proof>

lemma depth-diff-equiv0-left: p-equal p x 0  $\implies$  p-depth p (x - y) = p-depth p y
  <proof>

lemma depth-diff-equiv0-right: p-equal p y 0  $\implies$  p-depth p (x - y) = p-depth p x
  <proof>

lemma depth-diff-equiv:
  p-depth p (x - y) = p-depth p (w - z) if p-equal p x w and p-equal p y z
  <proof>

lemma depth-pre-nonarch-diff-left:
  p-depth p (x - y) = p-depth p x if  $\neg$  p-equal p x 0 and p-depth p x < p-depth p
  y
  <proof>

```

lemma *depth-pre-nonarch-diff-right*:

$p\text{-depth } p \ (x - y) = p\text{-depth } p \ y$ **if** $\neg p\text{-equal } p \ y \ 0$ **and** $p\text{-depth } p \ y < p\text{-depth } p \ x$
 ⟨proof⟩

lemma *depth-nonarch*:

$\llbracket \neg p\text{-equal } p \ x \ 0; \neg p\text{-equal } p \ y \ 0; \neg p\text{-equal } p \ (x + y) \ 0 \rrbracket$
 $\implies p\text{-depth } p \ (x + y) \geq \min \ (p\text{-depth } p \ x) \ (p\text{-depth } p \ y)$
 ⟨proof⟩

lemma *depth-nonarch-diff*:

$\llbracket \neg p\text{-equal } p \ x \ 0; \neg p\text{-equal } p \ y \ 0; \neg p\text{-equal } p \ (x - y) \ 0 \rrbracket$
 $\implies p\text{-depth } p \ (x - y) \geq \min \ (p\text{-depth } p \ x) \ (p\text{-depth } p \ y)$
 ⟨proof⟩

lemma *depth-sum-nonarch*:

$\text{finite } A \implies \neg p\text{-equal } p \ (\text{sum } f \ A) \ 0 \implies$
 $p\text{-depth } p \ (\text{sum } f \ A) \geq \text{Min } \{p\text{-depth } p \ (f \ a) \mid a. \ a \in A \wedge \neg p\text{-equal } p \ (f \ a) \ 0 \}$
 ⟨proof⟩

lemma *obtains-depth-sum-nonarch-witness*:

assumes $\text{finite } A$ **and** $\neg p\text{-equal } p \ (\text{sum } f \ A) \ 0$
obtains a
where $a \in A$ **and** $\neg p\text{-equal } p \ (f \ a) \ 0$ **and** $p\text{-depth } p \ (\text{sum } f \ A) \geq p\text{-depth } p \ (f \ a)$
 ⟨proof⟩

lemma *depth-add-cancel-imp-eq-depth*:

$p\text{-depth } p \ x = p\text{-depth } p \ y$ **if** $\neg p\text{-equal } p \ x \ 0$ **and** $p\text{-depth } p \ (x + y) > p\text{-depth } p \ x$
 ⟨proof⟩

lemma *depth-diff-cancel-imp-eq-depth-left*:

$p\text{-depth } p \ x = p\text{-depth } p \ y$ **if** $\neg p\text{-equal } p \ x \ 0$ **and** $p\text{-depth } p \ (x - y) > p\text{-depth } p \ x$
 ⟨proof⟩

lemma *depth-diff-cancel-imp-eq-depth-right*:

$p\text{-depth } p \ x = p\text{-depth } p \ y$ **if** $\neg p\text{-equal } p \ y \ 0$ **and** $p\text{-depth } p \ (x - y) > p\text{-depth } p \ y$
 ⟨proof⟩

lemma *level-closed-add*:

$p\text{-depth } p \ (x + y) \geq n$
if $\neg p\text{-equal } p \ (x + y) \ 0$ **and** $p\text{-depth } p \ x \geq n$ **and** $p\text{-depth } p \ y \geq n$
 ⟨proof⟩

lemma *level-closed-diff*:

$p\text{-depth } p \ (x - y) \geq n$

if $\neg p\text{-equal } p \ (x - y) \ 0$ **and** $p\text{-depth } p \ x \geq n$ **and** $p\text{-depth } p \ y \geq n$
 $\langle \text{proof} \rangle$

definition $p\text{-depth-set} :: 'a \Rightarrow \text{int} \Rightarrow 'b \text{ set}$
where $p\text{-depth-set } p \ n = \{x. \neg p\text{-equal } p \ x \ 0 \longrightarrow p\text{-depth } p \ x \geq n\}$

definition $\text{global-depth-set} :: \text{int} \Rightarrow 'b \text{ set}$
where
 $\text{global-depth-set } n =$
 $\{x. \forall p. \neg p\text{-equal } p \ x \ 0 \longrightarrow p\text{-depth } p \ x \geq n\}$

lemma global-depth-set : $\text{global-depth-set } n = (\bigcap p. p\text{-depth-set } p \ n)$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-set}D$:
 $\neg p\text{-equal } p \ x \ 0 \Longrightarrow x \in p\text{-depth-set } p \ n \Longrightarrow$
 $p\text{-depth } p \ x \geq n$
 $\langle \text{proof} \rangle$

lemma $\text{nonpos-}p\text{-depth-set}D$:
 $n \leq 0 \Longrightarrow x \in p\text{-depth-set } p \ n \Longrightarrow p\text{-depth } p \ x \geq n$
 $\langle \text{proof} \rangle$

lemma $\text{global-depth-set}D$:
 $\neg p\text{-equal } p \ x \ 0 \Longrightarrow x \in \text{global-depth-set } n \Longrightarrow$
 $p\text{-depth } p \ x \geq n$
 $\langle \text{proof} \rangle$

lemma $\text{nonpos-global-depth-set}D$:
 $n \leq 0 \Longrightarrow x \in \text{global-depth-set } n \Longrightarrow p\text{-depth } p \ x \geq n$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-set-}0$: $0 \in p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma $\text{global-depth-set-}0$: $0 \in \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-set-equiv}0$: $p\text{-equal } p \ x \ 0 \Longrightarrow x \in p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-set}I$:
 $x \in p\text{-depth-set } p \ n$ **if** $\neg p\text{-equal } p \ x \ 0 \longrightarrow p\text{-depth } p \ x \geq n$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-set-by-equiv}I$:
 $x \in p\text{-depth-set } p \ n$ **if** $p\text{-equal } p \ x \ y$ **and** $y \in p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *global-depth-setI*:

$x \in \text{global-depth-set } n$

if $\bigwedge p. \neg p\text{-equal } p \ x \ 0 \implies p\text{-depth } p \ x \geq n$

$\langle \text{proof} \rangle$

lemma *global-depth-setI'*: $x \in \text{global-depth-set } n$ **if** $\bigwedge p. p\text{-depth } p \ x \geq n$

$\langle \text{proof} \rangle$

lemma *global-depth-setI''*: $x \in \text{global-depth-set } n$ **if** $\bigwedge p. x \in p\text{-depth-set } p \ n$

$\langle \text{proof} \rangle$

lemma *p-depth-set-antimono*:

$n \leq m \implies p\text{-depth-set } p \ n \supseteq p\text{-depth-set } p \ m$

$\langle \text{proof} \rangle$

lemma *global-depth-set-antimono*:

$n \leq m \implies \text{global-depth-set } n \supseteq \text{global-depth-set } m$

$\langle \text{proof} \rangle$

lemma *p-depth-set-add*:

$x + y \in p\text{-depth-set } p \ n$ **if** $x \in p\text{-depth-set } p \ n$ **and** $y \in p\text{-depth-set } p \ n$

$\langle \text{proof} \rangle$

lemma *global-depth-set-add*:

$x + y \in \text{global-depth-set } n$ **if** $x \in \text{global-depth-set } n$ **and** $y \in \text{global-depth-set } n$

$\langle \text{proof} \rangle$

lemma *p-depth-set-uminus*: $-x \in p\text{-depth-set } p \ n$ **if** $x \in p\text{-depth-set } p \ n$

$\langle \text{proof} \rangle$

lemma *global-depth-set-uminus*: $-x \in \text{global-depth-set } n$ **if** $x \in \text{global-depth-set } n$

$\langle \text{proof} \rangle$

lemma *p-depth-set-minus*:

$x - y \in p\text{-depth-set } p \ n$ **if** $x \in p\text{-depth-set } p \ n$ **and** $y \in p\text{-depth-set } p \ n$

$\langle \text{proof} \rangle$

lemma *global-depth-set-minus*:

$x - y \in \text{global-depth-set } n$ **if** $x \in \text{global-depth-set } n$ **and** $y \in \text{global-depth-set } n$

$\langle \text{proof} \rangle$

lemma *p-depth-set-elt-set-plus*:

$x + o \ p\text{-depth-set } p \ n = p\text{-depth-set } p \ n$ **if** $x \in p\text{-depth-set } p \ n$

$\langle \text{proof} \rangle$

lemma *global-depth-set-elt-set-plus*:

$x + o \ \text{global-depth-set } n = \text{global-depth-set } n$ **if** $x \in \text{global-depth-set } n$

$\langle \text{proof} \rangle$

lemma *p-depth-set-elt-set-plus-closed*:
 $x + o \text{ } p\text{-depth-set } p \text{ } m \subseteq p\text{-depth-set } p \text{ } n$ **if** $m \geq n$ **and** $x \in p\text{-depth-set } p \text{ } n$
 ⟨proof⟩

lemma *global-depth-set-elt-set-plus-closed*:
 $x + o \text{ } global\text{-depth-set } m \subseteq global\text{-depth-set } n$
if $m \geq n$ **and** $x \in global\text{-depth-set } n$
 ⟨proof⟩

lemma *p-depth-set-plus-coeffs*:
 $set \text{ } xs \subseteq p\text{-depth-set } p \text{ } n \implies$
 $set \text{ } ys \subseteq p\text{-depth-set } p \text{ } n \implies$
 $set \text{ } (plus\text{-coeffs } xs \text{ } ys) \subseteq p\text{-depth-set } p \text{ } n$
 ⟨proof⟩

lemma *global-depth-set-plus-coeffs*:
 $set \text{ } (plus\text{-coeffs } xs \text{ } ys) \subseteq global\text{-depth-set } n$
if $set \text{ } xs \subseteq global\text{-depth-set } n$ **and** $set \text{ } ys \subseteq global\text{-depth-set } n$
 ⟨proof⟩

lemma *p-depth-set-poly-pCons*:
 $set \text{ } (coeffs \text{ } (pCons \text{ } a \text{ } f)) \subseteq p\text{-depth-set } p \text{ } n$
if $a \in p\text{-depth-set } p \text{ } n$ **and** $set \text{ } (coeffs \text{ } f) \subseteq p\text{-depth-set } p \text{ } n$
 ⟨proof⟩

lemma *p-depth-set-poly-add*:
 $set \text{ } (coeffs \text{ } (f + g)) \subseteq p\text{-depth-set } p \text{ } n$
if $set \text{ } (coeffs \text{ } f) \subseteq p\text{-depth-set } p \text{ } n$ **and** $set \text{ } (coeffs \text{ } g) \subseteq p\text{-depth-set } p \text{ } n$
 ⟨proof⟩

lemma *global-depth-set-poly-add*:
 $set \text{ } (coeffs \text{ } (f + g)) \subseteq global\text{-depth-set } n$
if $set \text{ } (coeffs \text{ } f) \subseteq global\text{-depth-set } n$
and $set \text{ } (coeffs \text{ } g) \subseteq global\text{-depth-set } n$
 ⟨proof⟩

end

locale *p-equality-depth-no-zero-divisors* =
p-equality-no-zero-divisors + *p-equality-depth*
 + **assumes** *depth-mult-additive*:
 $\neg p\text{-equal } p \text{ } (x * y) \text{ } 0 \implies p\text{-depth } p \text{ } (x * y) = p\text{-depth } p \text{ } x + p\text{-depth } p \text{ } y$
begin

lemma *depth-mult3-additive*:
 $p\text{-depth } p \text{ } (x * y * z) = p\text{-depth } p \text{ } x + p\text{-depth } p \text{ } y + p\text{-depth } p \text{ } z$
if $\neg p\text{-equal } p \text{ } (x * y * z) \text{ } 0$
 ⟨proof⟩

lemma *p-depth-set-times*:

$x * y \in p\text{-depth-set } p \ n$
if $n \geq 0$ **and** $x \in p\text{-depth-set } p \ n$ **and** $y \in p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-times*:

$x * y \in \text{global-depth-set } n$
if $n \geq 0$ **and** $x \in \text{global-depth-set } n$ **and** $y \in \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

lemma *poly-p-depth-set-poly*:

$\text{set } (\text{coeffs } f) \subseteq p\text{-depth-set } p \ n \implies \text{poly } f \ x \in p\text{-depth-set } p \ n$
if $n \geq 0$ **and** $x \in p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *poly-global-depth-set-poly*:

$\text{poly } f \ x \in \text{global-depth-set } n$
if $n \geq 0$
and $x \in \text{global-depth-set } n$
and $\text{set } (\text{coeffs } f) \subseteq \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-ideal*:

$x * y \in p\text{-depth-set } p \ n$
if $x \in p\text{-depth-set } p \ 0$ **and** $y \in p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-ideal*:

$x * y \in \text{global-depth-set } n$
if $x \in \text{global-depth-set } 0$ **and** $y \in \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

lemma *poly-p-depth-set-poly-ideal*:

$\text{set } (\text{coeffs } f) \subseteq p\text{-depth-set } p \ n \implies \text{poly } f \ x \in p\text{-depth-set } p \ n$
if $x \in p\text{-depth-set } p \ 0$
 $\langle \text{proof} \rangle$

lemma *poly-global-depth-set-poly-ideal*:

$\text{poly } f \ x \in \text{global-depth-set } n$
if $x \in \text{global-depth-set } 0$ **and** $\text{set } (\text{coeffs } f) \subseteq \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-poly-smult*:

$\text{set } (\text{coeffs } (\text{smult } x \ f)) \subseteq p\text{-depth-set } p \ n$
if $n \geq 0$ **and** $x \in p\text{-depth-set } p \ n$ **and** $\text{set } (\text{coeffs } f) \subseteq p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-poly-smult*:

```

    set (coeffs (smult x f))  $\subseteq$  global-depth-set n
  if n  $\geq$  0
  and x  $\in$  global-depth-set n
  and set (coeffs f)  $\subseteq$  global-depth-set n
  <proof>

lemma p-depth-set-poly-smult-ideal:
  set (coeffs (smult x f))  $\subseteq$  p-depth-set p n
  if x  $\in$  p-depth-set p 0 and set (coeffs f)  $\subseteq$  p-depth-set p n
  <proof>

lemma global-depth-set-poly-smult-ideal:
  set (coeffs (smult x f))  $\subseteq$  global-depth-set n
  if x  $\in$  global-depth-set 0 and set (coeffs f)  $\subseteq$  global-depth-set n
  <proof>

lemma p-depth-set-poly-times:
  set (coeffs f)  $\subseteq$  p-depth-set p n  $\implies$ 
    set (coeffs g)  $\subseteq$  p-depth-set p n  $\implies$ 
    set (coeffs (f * g))  $\subseteq$  p-depth-set p n
  if n  $\geq$  0
  <proof>

lemma global-depth-set-poly-times:
  set (coeffs (f * g))  $\subseteq$  global-depth-set n
  if n  $\geq$  0
  and set (coeffs g)  $\subseteq$  global-depth-set n
  and set (coeffs f)  $\subseteq$  global-depth-set n
  <proof>

lemma p-depth-set-poly-times-ideal:
  set (coeffs f)  $\subseteq$  p-depth-set p 0  $\implies$ 
    set (coeffs g)  $\subseteq$  p-depth-set p n  $\implies$ 
    set (coeffs (f * g))  $\subseteq$  p-depth-set p n
  <proof>

lemma global-depth-set-poly-times-ideal:
  set (coeffs (f * g))  $\subseteq$  global-depth-set n
  if set (coeffs f)  $\subseteq$  global-depth-set 0
  and set (coeffs g)  $\subseteq$  global-depth-set n
  <proof>

end

locale p-equality-depth-no-zero-divisors-1 =
  p-equality-no-zero-divisors-1 + p-equality-depth-no-zero-divisors
begin

```

lemma *depth-of-1[simp]*: $p\text{-depth } p \ 1 = 0$
 $\langle \text{proof} \rangle$

lemma *depth-of-neg1*: $p\text{-depth } p \ (-1) = 0$
 $\langle \text{proof} \rangle$

lemma *depth-pow-additive*: $p\text{-depth } p \ (x \wedge^n) = \text{int } n * p\text{-depth } p \ x$
 $\langle \text{proof} \rangle$

lemma *depth-prod-summative*:
 $\text{finite } X \implies \neg p\text{-equal } p \ (\prod X) \ 0 \implies$
 $p\text{-depth } p \ (\prod X) = \text{sum } (p\text{-depth } p) \ X$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-1*: $1 \in p\text{-depth-set } p \ 0$
 $\langle \text{proof} \rangle$

lemma *pos-p-depth-set-1*: $n > 0 \implies 1 \notin p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-1*: $1 \in \text{global-depth-set } 0$
 $\langle \text{proof} \rangle$

lemma *pos-global-depth-set-1*: $n > 0 \implies 1 \notin \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-neg1*: $-1 \in p\text{-depth-set } p \ 0$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-neg1*: $-1 \in \text{global-depth-set } 0$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-pow*:
 $x \wedge \text{Suc } k \in p\text{-depth-set } p \ n \text{ if } n \geq 0 \text{ and } x \in p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-pow*:
 $x \wedge \text{Suc } k \in \text{global-depth-set } n \text{ if } n \geq 0 \text{ and } x \in \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-0-pow*: $x \wedge k \in p\text{-depth-set } p \ 0 \text{ if } x \in p\text{-depth-set } p \ 0$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-0-pow*: $x \wedge k \in \text{global-depth-set } 0 \text{ if } x \in \text{global-depth-set } 0$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-of-nat*: $\text{of-nat } n \in p\text{-depth-set } p \ 0$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-of-nat*: of-nat $n \in \text{global-depth-set } 0$
 ⟨proof⟩

lemma *p-depth-set-of-int*: of-int $n \in \text{p-depth-set } p \ 0$
 ⟨proof⟩

lemma *global-depth-set-of-int*: of-int $n \in \text{global-depth-set } 0$
 ⟨proof⟩

lemma *p-depth-set-polyderiv*:
 set (coeffs f) $\subseteq \text{p-depth-set } p \ n \implies$
 set (coeffs (polyderiv f)) $\subseteq \text{p-depth-set } p \ n$
 ⟨proof⟩

lemma *global-depth-set-polyderiv*:
 set (coeffs (polyderiv f)) $\subseteq \text{global-depth-set } n$
 if set (coeffs f) $\subseteq \text{global-depth-set } n$
 ⟨proof⟩

lemma *poly-p-depth-set-polyderiv*:
 poly (polyderiv f) $x \in \text{p-depth-set } p \ n$
 if $n \geq 0$ and set (coeffs f) $\subseteq \text{p-depth-set } p \ n$ and $x \in \text{p-depth-set } p \ n$
 ⟨proof⟩

lemma *poly-global-depth-set-polyderiv*:
 poly (polyderiv f) $x \in \text{global-depth-set } n$
 if $n \geq 0$
 and set (coeffs f) $\subseteq \text{global-depth-set } n$
 and $x \in \text{global-depth-set } n$
 ⟨proof⟩

lemma *p-set-depth-additive-poly-poly*:
 set (coeffs f) $\subseteq \text{p-depth-set } p \ m \implies$
 set (coeffs (coeff (additive-poly-poly f) n)) $\subseteq \text{p-depth-set } p \ m$
 if $m \geq 0$
 ⟨proof⟩

lemma *global-depth-set-additive-poly-poly*:
 set (coeffs (coeff (additive-poly-poly f) n)) $\subseteq \text{global-depth-set } m$
 if $m \geq 0$ and set (coeffs f) $\subseteq \text{global-depth-set } m$
 ⟨proof⟩

lemma *poly-p-depth-set-additive-poly-poly*:
 poly (coeff (additive-poly-poly f) n) $x \in \text{p-depth-set } p \ m$
 if $m \geq 0$ and set (coeffs f) $\subseteq \text{p-depth-set } p \ m$ and $x \in \text{p-depth-set } p \ m$
 ⟨proof⟩

lemma *poly-global-depth-set-additive-poly-poly*:
 poly (coeff (additive-poly-poly f) n) $x \in \text{global-depth-set } m$

if $m \geq 0$
and $\text{set } (\text{coeffs } f) \subseteq \text{global-depth-set } m$
and $x \in \text{global-depth-set } m$
 $\langle \text{proof} \rangle$

lemma *sum-poly-additive-poly-poly-depth-bound*:
fixes $p :: 'a$ **and** $f :: 'b \text{ poly}$ **and** $x \ y :: 'b$ **and** $a \ b \ n :: \text{nat}$
defines $g \equiv \lambda i. \text{poly } (\text{coeff } (\text{additive-poly-poly } f) (\text{Suc } i)) \ x * y \wedge (\text{Suc } i)$
defines $S \equiv \text{sum } g \ \{a..b\}$
assumes $\text{coeffs}: \text{set } (\text{coeffs } f) \subseteq \text{p-depth-set } p \ 0$
and $\text{arg} : x \in \text{p-depth-set } p \ 0$
and $\text{sum} : \neg \text{p-equal } p \ S \ 0$
obtains j **where** $j \in \{a..b\}$ **and** $\text{p-depth } p \ S \geq \text{int } (\text{Suc } j) * \text{p-depth } p \ y$
 $\langle \text{proof} \rangle$

end

3.5.2 Depth of inverses and division

locale *p-equality-depth-div-inv* =
 $\text{p-equality-div-inv} + \text{p-equality-depth-no-zero-divisors-1}$
begin

lemma *inverse-depth*: $\text{p-depth } p \ (\text{inverse } x) = - \text{p-depth } p \ x$
 $\langle \text{proof} \rangle$

lemma *depth-powi-additive*: $\text{p-depth } p \ (\text{power-int } x \ n) = n * \text{p-depth } p \ x$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-inverse*: $\text{inverse } x \in \text{p-depth-set } p \ 0$ **if** $\text{p-depth } p \ x = 0$
 $\langle \text{proof} \rangle$

lemma *divide-depth*: $\text{p-depth } p \ (x / y) = \text{p-depth } p \ x - \text{p-depth } p \ y$ **if** $\neg \text{p-equal } p \ (x / y) \ 0$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-divide*:
 $x / y \in \text{p-depth-set } p \ n$ **if** $x \in \text{p-depth-set } p \ n$ **and** $\text{p-depth } p \ y = 0$
 $\langle \text{proof} \rangle$

lemma *p-depth-poly-deriv-quotient*:
 $\text{p-depth } p \ ((\text{poly } f \ a) / (\text{poly } (\text{polyderiv } f) \ a)) \geq n$
if $n \geq 0$
and $\text{set } (\text{coeffs } f) \subseteq \text{p-depth-set } p \ n$
and $a \in \text{p-depth-set } p \ n$
and $\neg \text{p-equal } p \ (\text{poly } f \ a) \ 0$
and $\neg \text{p-equal } p \ (\text{poly } (\text{polyderiv } f) \ a) \ 0$
and $\text{p-depth } p \ (\text{poly } f \ a) \geq 2 * \text{p-depth } p \ (\text{poly } (\text{polyderiv } f) \ a)$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-poly-deriv-quotient*:
 $(poly\ f\ a) / (poly\ (polyderiv\ f)\ a) \in p\text{-depth-set}\ p\ n$
if $n \geq 0$
and $set\ (coeffs\ f) \subseteq p\text{-depth-set}\ p\ n$
and $a \in p\text{-depth-set}\ p\ n$
and $\neg p\text{-equal}\ p\ (poly\ f\ a)\ 0$
and $\neg p\text{-equal}\ p\ (poly\ (polyderiv\ f)\ a)\ 0$
and $p\text{-depth}\ p\ (poly\ f\ a) \geq 2 * p\text{-depth}\ p\ (poly\ (polyderiv\ f)\ a)$
 $\langle proof \rangle$

end

3.6 Global equality and restriction of places

3.6.1 Locales

locale *global-p-equality* = *p-equality*
 $+$
fixes $p\text{-restrict} :: 'b::comm\text{-ring} \Rightarrow ('a \Rightarrow bool) \Rightarrow 'b$
assumes $p\text{-restrict-equiv} : P\ p \Longrightarrow p\text{-equal}\ p\ (p\text{-restrict}\ x\ P)\ x$
and $p\text{-restrict-equiv0} : \neg P\ p \Longrightarrow p\text{-equal}\ p\ (p\text{-restrict}\ x\ P)\ 0$
begin

lemma *p-restrict-equiv-sym*: $P\ p \Longrightarrow p\text{-equal}\ p\ x\ (p\text{-restrict}\ x\ P)$
 $\langle proof \rangle$

lemma *p-restrict-equiv0-iff*:
 $p\text{-equal}\ p\ (p\text{-restrict}\ x\ P)\ 0 \longleftrightarrow p\text{-equal}\ p\ x\ 0 \vee \neg P\ p$
 $\langle proof \rangle$

lemma *p-restrict-restrict-equiv-conj*:
 $p\text{-equal}\ p\ (p\text{-restrict}\ (p\text{-restrict}\ x\ P)\ Q)\ (p\text{-restrict}\ x\ (\lambda p. P\ p \wedge Q\ p))$
 $\langle proof \rangle$

lemma *p-restrict-restrict-equiv*: $p\text{-equal}\ p\ (p\text{-restrict}\ (p\text{-restrict}\ x\ P)\ P)\ (p\text{-restrict}\ x\ P)$
 $\langle proof \rangle$

lemma *p-restrict-0-equiv0*: $p\text{-equal}\ p\ (p\text{-restrict}\ 0\ P)\ 0$
 $\langle proof \rangle$

lemma *p-restrict-add-equiv*: $p\text{-equal}\ p\ (p\text{-restrict}\ (x + y)\ P)\ (p\text{-restrict}\ x\ P + p\text{-restrict}\ y\ P)$
 $\langle proof \rangle$

lemma *p-restrict-add-mixed-equiv*:
 $p\text{-equal}\ p\ (p\text{-restrict}\ x\ P + p\text{-restrict}\ y\ Q)\ (x + y)$ **if** $P\ p$ **and** $Q\ p$
 $\langle proof \rangle$

lemma *p-restrict-add-mixed-equiv-drop-right:*

p-equal p (p-restrict x P + p-restrict y Q) x if P p and $\neg$ Q p

<proof>

lemma *p-restrict-add-mixed-equiv-drop-left:*

p-equal p (p-restrict x P + p-restrict y Q) y if $\neg$ P p and Q p

<proof>

lemma *p-restrict-add-mixed-equiv-drop-both:*

p-equal p (p-restrict x P + p-restrict y Q) 0 if $\neg$ P p and $\neg$ Q p

<proof>

lemma *p-restrict-uminus-equiv:* *p-equal p (p-restrict $(-x)$ P) $(- p-restrict x P)$*

<proof>

lemma *p-restrict-minus-equiv:* *p-equal p (p-restrict $(x - y)$ P) $(p-restrict x P - p-restrict y P)$*

<proof>

lemma *p-restrict-minus-mixed-equiv:*

p-equal p (p-restrict x P - p-restrict y Q) $(x - y)$ if P p and Q p

<proof>

lemma *p-restrict-minus-mixed-equiv-drop-right:*

p-equal p (p-restrict x P - p-restrict y Q) x if P p and $\neg$ Q p

<proof>

lemma *p-restrict-minus-mixed-equiv-drop-left:*

p-equal p (p-restrict x P - p-restrict y Q) $(-y)$ if $\neg$ P p and Q p

<proof>

lemma *p-restrict-minus-mixed-equiv-drop-both:*

p-equal p (p-restrict x P - p-restrict y Q) 0 if $\neg$ P p and $\neg$ Q p

<proof>

lemma *p-restrict-times-equiv:*

*p-equal p ((p-restrict x P) * (p-restrict y Q))*

*(p-restrict $(x * y)$ $(\lambda p. P p \wedge Q p)$)*

<proof>

lemma *p-restrict-times-equiv':*

*p-equal p ((p-restrict x P) * (p-restrict y P)) $(p-restrict (x * y) P)$*

<proof>

lemma *p-restrict-p-equalI:*

p-equal p (p-restrict x P) y

if *P p $\longrightarrow$ p-equal p x y*

and *$\neg$ P p $\longrightarrow$ p-equal p y 0*

<proof>

lemma *p-restrict-p-equal-p-restrictI*:
 $p\text{-equal } p \text{ (p-restrict } x \text{ } P) \text{ (p-restrict } y \text{ } Q)$
if $(P \text{ } p \wedge Q \text{ } p) \longrightarrow p\text{-equal } p \text{ } x \text{ } y$
and $(P \text{ } p \wedge \neg Q \text{ } p) \longrightarrow p\text{-equal } p \text{ } x \text{ } 0$
and $(\neg P \text{ } p \wedge Q \text{ } p) \longrightarrow p\text{-equal } p \text{ } y \text{ } 0$
 $\langle \text{proof} \rangle$

lemma *p-restrict-p-equal-p-restrict-by-sameI*:
 $p\text{-equal } p \text{ (p-restrict } x \text{ } P) \text{ (p-restrict } y \text{ } P) \text{ if } P \text{ } p \longrightarrow p\text{-equal } p \text{ } x \text{ } y$
 $\langle \text{proof} \rangle$

end

locale *global-p-equality-div-inv* =
 $p\text{-equality-div-inv } p\text{-equal} + \text{global-p-equality } p\text{-equal } p\text{-restrict}$
for $p\text{-equal} \quad ::$
 $'a \Rightarrow 'b :: \{\text{comm-ring-1, inverse, divide-trivial}\} \Rightarrow$
 $'b \Rightarrow \text{bool}$
and $p\text{-restrict} :: 'b \Rightarrow ('a \Rightarrow \text{bool}) \Rightarrow 'b$
begin

lemma *inverse-p-restrict-equiv*: $p\text{-equal } p \text{ (inverse (p-restrict } x \text{ } P)) \text{ (p-restrict (inverse } x \text{) } P)$
 $\langle \text{proof} \rangle$

end

locale *global-p-equal* = *global-p-equality*
 $+ \text{assumes } \text{global-imp-eq}: \text{globally-p-equal } x \text{ } y \Longrightarrow x = y$
begin

lemma *global-eq-iff*: $\text{globally-p-equal } x \text{ } y \longleftrightarrow x = y$
 $\langle \text{proof} \rangle$

lemma *p-restrict-true[simp]*: $p\text{-restrict } x \text{ } (\lambda x. \text{True}) = x$
 $\langle \text{proof} \rangle$

lemma *p-restrict-false[simp]*: $p\text{-restrict } x \text{ } (\lambda x. \text{False}) = 0$
 $\langle \text{proof} \rangle$

lemma *p-restrict-restrict*:
 $p\text{-restrict (p-restrict } x \text{ } P) \text{ } Q = p\text{-restrict } x \text{ } (\lambda p. P \text{ } p \wedge Q \text{ } p)$
 $\langle \text{proof} \rangle$

lemma *p-restrict-restrict'[simp]*: $p\text{-restrict (p-restrict } x \text{ } P) \text{ } P = p\text{-restrict } x \text{ } P$
 $\langle \text{proof} \rangle$

lemma *p-restrict-image-restrict*: $p\text{-restrict } x \ P = x \text{ if } x \in (\lambda z. p\text{-restrict } z \ P) \text{ ' } B$
 $\langle \text{proof} \rangle$

lemma *p-restrict-zero*: $p\text{-restrict } 0 \ P = 0$
 $\langle \text{proof} \rangle$

lemma *p-restrict-add*: $p\text{-restrict } (x + y) \ P = p\text{-restrict } x \ P + p\text{-restrict } y \ P$
 $\langle \text{proof} \rangle$

lemma *p-restrict-uminus*: $p\text{-restrict } (-x) \ P = - \ p\text{-restrict } x \ P$
 $\langle \text{proof} \rangle$

lemma *p-restrict-minus*: $p\text{-restrict } (x - y) \ P = p\text{-restrict } x \ P - p\text{-restrict } y \ P$
 $\langle \text{proof} \rangle$

lemma *p-restrict-times*:
 $(p\text{-restrict } x \ P) * (p\text{-restrict } y \ Q) = p\text{-restrict } (x * y) \ (\lambda p. P \ p \wedge Q \ p)$
 $\langle \text{proof} \rangle$

lemma *times-p-restrict*: $x * (p\text{-restrict } y \ P) = p\text{-restrict } (x * y) \ P$
 $\langle \text{proof} \rangle$

lemmas *p-restrict-pre-finite-adele-simps* =
 $p\text{-restrict-zero } p\text{-restrict-add } p\text{-restrict-uminus } p\text{-restrict-minus } p\text{-restrict-times}$

lemma *p-equal-0-iff-p-restrict-eq-0*: $p\text{-equal } p \ x \ 0 \longleftrightarrow p\text{-restrict } x \ ((=) \ p) = 0$
 $\langle \text{proof} \rangle$

lemma *p-equal-iff-p-restrict-eq*:
 $p\text{-equal } p \ x \ y \longleftrightarrow p\text{-restrict } x \ ((=) \ p) = p\text{-restrict } y \ ((=) \ p)$
 $\langle \text{proof} \rangle$

lemma *p-restrict-eqI*:
 $y = p\text{-restrict } x \ P$
if $P: \bigwedge p. P \ p \implies p\text{-equal } p \ y \ x$
and $\text{not-}P: \bigwedge p. \neg P \ p \implies p\text{-equal } p \ y \ 0$
 $\langle \text{proof} \rangle$

lemma *p-restrict-eq-p-restrict-mixedI*:
 $p\text{-restrict } x \ P = p\text{-restrict } y \ Q$
if $\text{both: } \bigwedge p::'a. P \ p \implies Q \ p \implies p\text{-equal } p \ x \ y$
and $P : \bigwedge p::'a. P \ p \implies \neg Q \ p \implies p\text{-equal } p \ x \ 0$
and $Q : \bigwedge p::'a. \neg P \ p \implies Q \ p \implies p\text{-equal } p \ y \ 0$
 $\langle \text{proof} \rangle$

lemma *p-restrict-eq-p-restrictI*:
 $p\text{-restrict } x \ P = p\text{-restrict } y \ P \text{ if } \bigwedge p::'a. P \ p \implies p\text{-equal } p \ x \ y$
 $\langle \text{proof} \rangle$

lemma *p-restrict-eq-p-restrict-iff*:

$p\text{-restrict } x \ P = p\text{-restrict } y \ P \longleftrightarrow$
 $(\forall p::'a. \ P \ p \longrightarrow p\text{-equal } p \ x \ y)$
 $\langle \text{proof} \rangle$

lemma *p-restrict-eq-zero-iff*:

$p\text{-restrict } x \ P = 0 \longleftrightarrow (\forall p. \ P \ p \longrightarrow p\text{-equal } p \ x \ 0)$
 $\langle \text{proof} \rangle$

lemma *p-restrict-decomp*:

$x = p\text{-restrict } x \ P + p\text{-restrict } x \ (\lambda p. \ \neg \ P \ p)$
 $\langle \text{proof} \rangle$

lemma *restrict-complement*:

$\text{range } (\lambda x. \ p\text{-restrict } x \ P) + \text{range } (\lambda x. \ p\text{-restrict } x \ (\lambda p. \ \neg \ P \ p))$
 $= \text{UNIV}$
 $\langle \text{proof} \rangle$

end

locale *global-p-equal-no-zero-divisors* =

global-p-equal *p-equal* *p-restrict* + *p-equality-no-zero-divisors* *p-equal*
for *p-equal* :: $'a \Rightarrow 'b::\text{comm-ring} \Rightarrow 'b \Rightarrow \text{bool}$
and *p-restrict* :: $'b \Rightarrow ('a \Rightarrow \text{bool}) \Rightarrow 'b$

locale *global-p-equal-no-zero-divisors-1* =

global-p-equal-no-zero-divisors *p-equal* *p-restrict* + *p-equality-no-zero-divisors-1*
p-equal
for *p-equal* :: $'a \Rightarrow 'b::\text{comm-ring-1} \Rightarrow 'b \Rightarrow \text{bool}$
and *p-restrict* :: $'b \Rightarrow ('a \Rightarrow \text{bool}) \Rightarrow 'b$
begin

lemma *pow-eq-0-iff*: $x \wedge^n = 0 \longleftrightarrow x = 0 \wedge n \neq 0$ **for** $x :: 'b$

$\langle \text{proof} \rangle$

end

locale *global-p-equal-div-inv* =

global-p-equality-div-inv *p-equal* *p-restrict*
+ *global-p-equal* *p-equal* *p-restrict*
for *p-equal* ::
 $'a \Rightarrow 'b::\{\text{comm-ring-1}, \text{inverse}, \text{divide-trivial}\} \Rightarrow$
 $'b \Rightarrow \text{bool}$
and *p-restrict* :: $'b \Rightarrow ('a \Rightarrow \text{bool}) \Rightarrow 'b$
begin

lemma *power-int-eq-0-iff*:

$\text{power-int } x \ n = 0 \longleftrightarrow x = 0 \wedge n \neq 0$ **for** $x :: 'b$
 $\langle \text{proof} \rangle$

lemma *inverse-eq0-iff*: $\text{inverse } x = 0 \longleftrightarrow x = 0$ **for** $x :: 'b$

$\langle \text{proof} \rangle$

lemma *inverse-neg-1* [simp]: $\text{inverse } (-1 :: 'b) = -1$

$\langle \text{proof} \rangle$

lemma *inverse-uminus*: $\text{inverse } (-x) = - \text{inverse } x$ **for** $x :: 'b$

$\langle \text{proof} \rangle$

lemma *global-right-inverse*: $x * \text{inverse } x = 1$ **if** $\forall p. \neg p\text{-equal } p \ x \ 0$

$\langle \text{proof} \rangle$

lemma *right-inverse*: $x * \text{inverse } x = p\text{-restrict } 1 \ (\lambda p. \neg p\text{-equal } p \ x \ 0)$

$\langle \text{proof} \rangle$

lemma *global-left-inverse*: $\text{inverse } x * x = 1$ **if** $\forall p. \neg p\text{-equal } p \ x \ 0$

$\langle \text{proof} \rangle$

lemma *left-inverse*: $\text{inverse } x * x = p\text{-restrict } 1 \ (\lambda p. \neg p\text{-equal } p \ x \ 0)$

$\langle \text{proof} \rangle$

lemma *inverse-unique*: $\text{inverse } x = y$ **if** $x * y = 1$ **for** $x \ y :: 'b$

$\langle \text{proof} \rangle$

lemma *inverse-p-restrict*: $\text{inverse } (p\text{-restrict } x \ P) = p\text{-restrict } (\text{inverse } x) \ P$

$\langle \text{proof} \rangle$

lemma *power-diff-conv-inverse*:

$x \wedge (n - m) = x \wedge n * \text{inverse } x \wedge m$ **if** $m < n$
for $x :: 'b$

$\langle \text{proof} \rangle$

lemma *power-int-add-1*: $\text{power-int } x \ (m + 1) = \text{power-int } x \ m * x$ **if** $m \neq -1$ **for**

$x :: 'b$

$\langle \text{proof} \rangle$

lemma *power-int-add*:

$\text{power-int } x \ (m + n) = \text{power-int } x \ m * \text{power-int } x \ n$

if $(\forall p. \neg p\text{-equal } p \ x \ 0) \vee m + n \neq 0$ **for** $x :: 'b$

$\langle \text{proof} \rangle$

lemma *power-int-diff*:

$\text{power-int } x \ (m - n) = \text{power-int } x \ m / \text{power-int } x \ n$ **if** $m \neq n$ **for** $x :: 'b$

$\langle \text{proof} \rangle$

lemma *power-int-minus-mult*: $\text{power-int } x \ (n - 1) * x = \text{power-int } x \ n$ **if** $n \neq 0$
for $x :: 'b$
 $\langle \text{proof} \rangle$

lemma *power-int-not-eq-zero*:
 $x \neq 0 \vee n = 0 \implies \text{power-int } x \ n \neq 0$ **for** $x :: 'b$
 $\langle \text{proof} \rangle$

lemma *minus-divide-right*: $a / (-b) = -(a / b)$ **for** $a \ b :: 'b$
 $\langle \text{proof} \rangle$

lemma *minus-divide-divide*: $(-a) / (-b) = a / b$ **for** $a \ b :: 'b$
 $\langle \text{proof} \rangle$

lemma *divide-minus1* [simp]: $x / (-1) = -x$ **for** $x :: 'b$
 $\langle \text{proof} \rangle$

lemma *divide-self*: $x / x = p\text{-restrict } 1 \ (\lambda p. \neg p\text{-equal } p \ x \ 0)$
 $\langle \text{proof} \rangle$

lemma *global-divide-self*: $x / x = 1$ **if** $\forall p. \neg p\text{-equal } p \ x \ 0$
 $\langle \text{proof} \rangle$

lemma *global-mult-divide-mult-cancel-right*:
 $(a * b) / (c * b) = a / c$ **if** $\forall p. \neg p\text{-equal } p \ b \ 0$
 $\langle \text{proof} \rangle$

lemma *global-mult-divide-mult-cancel-left*:
 $(c * a) / (c * b) = a / b$ **if** $\forall p. \neg p\text{-equal } p \ c \ 0$
 $\langle \text{proof} \rangle$

lemma *divide-p-restrict-left*: $(p\text{-restrict } x \ P) / y = p\text{-restrict } (x / y) \ P$
 $\langle \text{proof} \rangle$

lemma *divide-p-restrict-right*: $x / (p\text{-restrict } y \ P) = p\text{-restrict } (x / y) \ P$
 $\langle \text{proof} \rangle$

lemma *inj-divide-right*:
 $\text{inj } (\lambda b. b / a) \longleftrightarrow (\forall p. \neg p\text{-equal } p \ a \ 0)$
 $\langle \text{proof} \rangle$

end

locale *global-p-equality-depth* = *global-p-equality* + *p-equality-depth*
begin

lemma *globally-p-equal-p-depth*: $\text{globally-p-equal } x \ y \implies p\text{-depth } p \ x = p\text{-depth } p \ y$

y
 $\langle \text{proof} \rangle$

lemma *p-restrict-depth*:

$P \ p \implies p\text{-depth } p \ (p\text{-restrict } x \ P) = p\text{-depth } p \ x$
 $\neg P \ p \implies p\text{-depth } p \ (p\text{-restrict } x \ P) = 0$
 $\langle \text{proof} \rangle$

lemma *p-restrict-add-mixed-depth-drop-right*:

$p\text{-depth } p \ (p\text{-restrict } x \ P + p\text{-restrict } y \ Q) = p\text{-depth } p \ x$ **if** $P \ p$ **and** $\neg Q \ p$
 $\langle \text{proof} \rangle$

lemma *p-restrict-add-mixed-depth-drop-left*:

$p\text{-depth } p \ (p\text{-restrict } x \ P + p\text{-restrict } y \ Q) = p\text{-depth } p \ y$ **if** $\neg P \ p$ **and** $Q \ p$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-p-restrict*:

$p\text{-restrict } x \ P \in p\text{-depth-set } p \ n$ **if** $P \ p \implies x \in p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-p-restrict*:

$p\text{-restrict } x \ P \in \text{global-depth-set } n$ **if** $x \in \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-closed-under-p-restrict-image*:

$(\lambda x. p\text{-restrict } x \ P) \ ' p\text{-depth-set } p \ n \subseteq p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-closed-under-p-restrict-image*:

$(\lambda x. p\text{-restrict } x \ P) \ ' \text{global-depth-set } n \subseteq \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-p-restrictI*:

$p\text{-restrict } x \ P \in \text{global-depth-set } n$
if $\bigwedge p. P \ p \implies x \in p\text{-depth-set } p \ n$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-image-under-one-p-restrict-image*:

$(\lambda x. p\text{-restrict } x \ ((=) \ p)) \ ' p\text{-depth-set } p \ n \subseteq \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

end

locale *global-p-equality-depth-no-zero-divisors-1* =

$p\text{-equality-depth-no-zero-divisors-1 } p\text{-equal } p\text{-depth}$
 $+ \text{global-p-equality-depth } p\text{-equal } p\text{-restrict } p\text{-depth}$
for $p\text{-equal} \quad \therefore$
 $\ 'a \Rightarrow 'b::\text{comm-ring-1} \Rightarrow 'b \Rightarrow \text{bool}$

```

and p-restrict :: 'b  $\Rightarrow$  ('a  $\Rightarrow$  bool)  $\Rightarrow$  'b
and p-depth    :: 'a  $\Rightarrow$  'b  $\Rightarrow$  int

locale global-p-equality-depth-div-inv =
  p-equality-depth-div-inv + global-p-equality-depth
begin

sublocale global-p-equality-div-inv <proof>

lemma global-depth-set-inverse:
  p-restrict (inverse x) ( $\lambda p. p\text{-depth } p \ x = 0$ )  $\in$  global-depth-set 0
  <proof>

lemma global-depth-set-divide:
  p-restrict (x / y) ( $\lambda p. p\text{-depth } p \ y = 0$ )  $\in$  global-depth-set n
  if x  $\in$  global-depth-set n
  <proof>

lemma global-depth-set-divideI:
  p-restrict (x / y) ( $\lambda p. p\text{-depth } p \ y = 0$ )  $\in$  global-depth-set n
  if  $\bigwedge p. p\text{-depth } p \ y = 0 \implies x \in p\text{-depth-set } p \ n$ 
  <proof>

end

locale global-p-equal-depth = global-p-equal + global-p-equality-depth
begin

lemma p-depth-set-vs-global-depth-set-image-under-one-p-restrict-image:
  ( $\lambda x. p\text{-restrict } x \ ((=) \ p)$ ) ' p-depth-set p n =
  ( $\lambda x. p\text{-restrict } x \ ((=) \ p)$ ) ' global-depth-set n
  <proof>

end

locale global-p-equal-depth-no-zero-divisors =
  global-p-equal-no-zero-divisors + p-equality-depth-no-zero-divisors
begin
sublocale global-p-equal-depth <proof>
end

locale global-p-equal-depth-div-inv =
  global-p-equality-depth-div-inv + global-p-equal-depth-no-zero-divisors
begin

sublocale global-p-equal-div-inv <proof>

```

lemma *p-depth-set-eq-coset*:

*p-depth-set p n = (w powi n) *o (p-depth-set p 0)*
if *p-depth p w = 1* **and** $\forall p. \neg p\text{-equal } p \ w \ 0$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-eq-coset*:

*global-depth-set n = (w powi n) *o (global-depth-set 0)* **if** $\forall p. p\text{-depth } p \ w = 1$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-eq-coset'*:

*p-depth-set p 0 = (w powi (-n)) *o (p-depth-set p n)*
if *p-depth p w = 1* **and** $\forall p. \neg p\text{-equal } p \ w \ 0$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-eq-coset'*:

*global-depth-set 0 = (w powi (-n)) *o (global-depth-set n)* **if** $\forall p. p\text{-depth } p \ w = 1$
 $\langle \text{proof} \rangle$

end

3.6.2 Embedding of base type constants

locale *global-p-equality-w-consts = global-p-equality p-equal p-restrict*

for *p-equal* $::$

'a $\Rightarrow$ *'b::comm-ring-1* $\Rightarrow$ *'b* $\Rightarrow$ *bool*

and *p-restrict* $::$ *'b* $\Rightarrow$ (*'a* $\Rightarrow$ *bool*) $\Rightarrow$ *'b*

+

fixes *func-of-consts* $::$ (*'a* $\Rightarrow$ *'c::ring-1*) $\Rightarrow$ *'b*

assumes *consts-1* $[simp]:$ *func-of-consts 1 = 1*

and *consts-diff* $[simp]:$ *func-of-consts (f - g) = func-of-consts f - func-of-consts g*

and *consts-mult* $[simp]:$

*func-of-consts (f * g) = func-of-consts f * func-of-consts g*

begin

abbreviation *consts* $\equiv$ *func-of-consts*

abbreviation *const a* $\equiv$ *consts* ($\lambda p. a$)

lemma *consts-0[simp]:* *consts 0 = 0*

$\langle \text{proof} \rangle$

lemma *const-0[simp]:* *const 0 = 0*

$\langle \text{proof} \rangle$

lemma *const-1[simp]:* *const 1 = 1*

$\langle \text{proof} \rangle$

lemma *consts-uminus* [*simp*]: *consts* $(- f) = - \text{consts } f$
 ⟨*proof*⟩

lemma *const-uminus* [*simp*]: *const* $(-a) = - \text{const } a$
 ⟨*proof*⟩

lemma *const-diff* [*simp*]: *const* $(a - b) = \text{const } a - \text{const } b$
 ⟨*proof*⟩

lemma *consts-add* [*simp*]: *consts* $(f + g) = \text{consts } f + \text{consts } g$
 ⟨*proof*⟩

lemma *const-add* [*simp*]: *const* $(a + b) = \text{const } a + \text{const } b$
 ⟨*proof*⟩

lemma *const-mult* [*simp*]: *const* $(a * b) = \text{const } a * \text{const } b$
 ⟨*proof*⟩

lemma *const-of-nat*: *const* $(\text{of-nat } n) = \text{of-nat } n$
 ⟨*proof*⟩

lemma *const-of-int*: *const* $(\text{of-int } z) = \text{of-int } z$
 ⟨*proof*⟩

end

locale *global-p-equal-w-consts* =
global-p-equal p-equal p-restrict
 + *global-p-equality-w-consts p-equal p-restrict func-of-consts*
for *p-equal* ::
 'a $\Rightarrow$ *'b::comm-ring-1* $\Rightarrow$ *'b* $\Rightarrow$ *bool*
and *p-restrict* :: *'b* $\Rightarrow$ (*'a* $\Rightarrow$ *bool*) $\Rightarrow$ *'b*
and *func-of-consts* :: (*'a* $\Rightarrow$ *'c::ring-1*) $\Rightarrow$ *'b*

locale *global-p-equality-w-inj-consts* = *global-p-equality-w-consts*
 +
assumes *p-equal-func-of-consts-0*: *p-equal* *p* $(\text{consts } f) \ 0 \longleftrightarrow f \ p = 0$
begin

lemmas *p-equal-func-of-const-0* = *p-equal-func-of-consts-0*[*of* - $\lambda p. \ -$]

lemma *p-equal-func-of-consts*:
p-equal *p* $(\text{consts } f) (\text{consts } g) \longleftrightarrow f \ p = g \ p$
 ⟨*proof*⟩

lemma *globally-p-equal-func-of-consts*:
globally-p-equal $(\text{func-of-consts } f) (\text{func-of-consts } g) \longleftrightarrow$

```

    (∀ p. f p = g p)
  ⟨proof⟩

lemma const-p-equal: p-equal p (const a) (const b) ⇒ a = b
  ⟨proof⟩

end

locale global-p-equality-w-inj-consts-char-0 =
  global-p-equality-w-inj-consts p-equal p-restrict func-of-consts
  for p-equal      ::
    'a ⇒ 'b::comm-ring-1 ⇒ 'b ⇒ bool
  and p-restrict  :: 'b ⇒ ('a ⇒ bool) ⇒ 'b
  and func-of-consts :: ('a ⇒ 'c::ring-char-0) ⇒ 'b
begin

lemma const-of-nat-p-equal-0-iff: p-equal p (of-nat n) 0 ⇔ n = 0
  ⟨proof⟩

lemma const-of-int-p-equal-0-iff: p-equal p (of-int z) 0 ⇔ z = 0
  ⟨proof⟩

end

locale global-p-equality-div-inv-w-inj-consts =
  p-equality-div-inv + global-p-equality-w-inj-consts
begin

lemma consts-divide-self-equiv: p-equal p ((consts f) / (consts f)) 1 if f p ≠ 0
  ⟨proof⟩

lemma const-divide-self-equiv: p-equal p ((const c) / (const c)) 1 if c ≠ 0
  ⟨proof⟩

end

locale global-p-equal-w-inj-consts =
  global-p-equal p-equal p-restrict
+ global-p-equality-w-inj-consts p-equal p-restrict
  for p-equal      ::
    'a ⇒ 'b::comm-ring-1 ⇒ 'b ⇒ bool
  and p-restrict  :: 'b ⇒ ('a ⇒ bool) ⇒ 'b
begin

lemma consts-eq-iff: consts f = consts g ⇔ f = g
  ⟨proof⟩

```

lemma *consts-eq-0-iff*: $\text{consts } f = 0 \longleftrightarrow f = 0$
 $\langle \text{proof} \rangle$

lemma *inj-consts*: $\text{inj } \text{consts}$
 $\langle \text{proof} \rangle$

lemma *const-eq-iff*: $\text{const } a = \text{const } b \longleftrightarrow a = b$
 $\langle \text{proof} \rangle$

lemma *const-eq-0-iff*: $\text{const } a = 0 \longleftrightarrow a = 0$
 $\langle \text{proof} \rangle$

lemma *inj-const*: $\text{inj } \text{const}$
 $\langle \text{proof} \rangle$

end

locale *global-p-equal-div-inv-w-inj-consts* =
 $\text{global-p-equal-div-inv} + \text{global-p-equal-w-inj-consts}$
begin

sublocale *global-p-equality-div-inv-w-inj-consts* $\langle \text{proof} \rangle$

lemma *consts-divide-self*:
 $(\text{consts } f) / (\text{consts } f) = p\text{-restrict } 1 \ (\lambda p::'a. f \ p \neq 0)$
 $\langle \text{proof} \rangle$

lemma *const-right-inverse* [simp]: $\text{const } a * \text{inverse } (\text{const } a) = 1 \text{ if } a \neq 0$
 $\langle \text{proof} \rangle$

lemma *const-left-inverse* [simp]: $\text{inverse } (\text{const } a) * \text{const } a = 1 \text{ if } a \neq 0$
 $\langle \text{proof} \rangle$

lemma *const-divide-self*: $(\text{const } c) / (\text{const } c) = 1 \text{ if } c \neq 0$
 $\langle \text{proof} \rangle$

lemma *mult-const-divide-mult-const-cancel-right*:
 $(y * \text{const } a) / (z * \text{const } a) = y / z \text{ if } a \neq 0$
 $\langle \text{proof} \rangle$

lemma *mult-const-divide-mult-const-cancel-left*:
 $(\text{const } a * y) / (\text{const } a * z) = y / z \text{ if } a \neq 0$
 $\langle \text{proof} \rangle$

lemma *mult-const-divide-mult-const-cancel-left-right*:
 $(\text{const } a * y) / (z * \text{const } a) = y / z \text{ if } a \neq 0$
 $\langle \text{proof} \rangle$

```

lemma mult-const-divide-mult-const-cancel-right-left:
   $(y * \text{const } a) / (\text{const } a * z) = y / z$  if  $a \neq 0$ 
   $\langle \text{proof} \rangle$ 

end

locale global-p-equal-div-inv-w-inj-consts-char-0 =
  global-p-equality-w-inj-consts-char-0 p-equal p-restrict func-of-consts
+ global-p-equal-div-inv-w-inj-consts p-equal p-restrict func-of-consts
  for p-equal ::
     $'a \Rightarrow 'b :: \{ \text{comm-ring-1}, \text{ring-char-0}, \text{inverse}, \text{divide-trivial} \} \Rightarrow$ 
     $'b \Rightarrow \text{bool}$ 
  and p-restrict ::  $'b \Rightarrow ('a \Rightarrow \text{bool}) \Rightarrow 'b$ 
  and func-of-consts ::  $('a \Rightarrow 'c :: \text{ring-char-0}) \Rightarrow 'b$ 
begin

context
  fixes  $n :: \text{nat}$ 
  assumes  $n \neq 0$ 
begin

lemma left-inverse-const-of-nat[simp]:  $\text{inverse } (\text{of-nat } n :: 'b) * \text{of-nat } n = 1$ 
   $\langle \text{proof} \rangle$ 

lemma right-inverse-const-of-nat[simp]:  $\text{of-nat } n * \text{inverse } (\text{of-nat } n :: 'b) = 1$ 
   $\langle \text{proof} \rangle$ 

lemma divide-self-const-of-nat[simp]:  $(\text{of-nat } n :: 'b) / \text{of-nat } n = 1$ 
   $\langle \text{proof} \rangle$ 

lemma mult-const-of-nat-divide-const-mult-of-nat-cancel-right:
   $(a * \text{of-nat } n) / (b * \text{of-nat } n) = a / b$  for  $a \ b :: 'b$ 
   $\langle \text{proof} \rangle$ 

lemma mult-const-of-nat-divide-const-mult-of-nat-cancel-left:
   $(\text{of-nat } n * a) / (\text{of-nat } n * b) = a / b$  for  $a \ b :: 'b$ 
   $\langle \text{proof} \rangle$ 

end

context
  fixes  $z :: \text{int}$ 
  assumes  $z \neq 0$ 
begin

lemma left-inverse-const-of-int[simp]:  $\text{inverse } (\text{of-int } z :: 'b) * \text{of-int } z = 1$ 
   $\langle \text{proof} \rangle$ 

```

lemma *right-inverse-const-of-int[simp]*: $\text{of-int } z * \text{inverse } (\text{of-int } z :: 'b) = 1$
 $\langle \text{proof} \rangle$

lemma *divide-self-const-of-int[simp]*: $(\text{of-int } z :: 'b) / \text{of-int } z = 1$
 $\langle \text{proof} \rangle$

lemma *mult-const-of-int-divide-const-mult-of-int-cancel-right*:
 $(a * \text{of-int } z) / (b * \text{of-int } z) = a / b$ **for** $a \ b :: 'b$
 $\langle \text{proof} \rangle$

lemma *mult-const-of-int-divide-const-mult-of-int-cancel-left*:
 $(\text{of-int } z * a) / (\text{of-int } z * b) = a / b$ **for** $a \ b :: 'b$
 $\langle \text{proof} \rangle$

end

context
includes *rat.lifting*
begin

lift-definition *const-of-rat* :: $\text{rat} \Rightarrow 'b$
is $\lambda x. \text{of-int } (\text{fst } x) / \text{of-int } (\text{snd } x)$
 $\langle \text{proof} \rangle$

lemma *inj-const-of-rat*: $\text{inj } \text{const-of-rat}$
 $\langle \text{proof} \rangle$

lemma *const-of-rat-rat*: $\text{const-of-rat } (\text{Rat.Fract } a \ b) = (\text{of-int } a :: 'b) / \text{of-int } b$
 $\langle \text{proof} \rangle$

lemma *const-of-rat-0 [simp]*: $\text{const-of-rat } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *const-of-rat-1 [simp]*: $\text{const-of-rat } 1 = 1$
 $\langle \text{proof} \rangle$

lemma *const-of-rat-minus*: $\text{const-of-rat } (- a) = - \text{const-of-rat } a$
 $\langle \text{proof} \rangle$

lemma *const-of-rat-add*: $\text{const-of-rat } (a + b) = \text{const-of-rat } a + \text{const-of-rat } b$
 $\langle \text{proof} \rangle$

lemma *const-of-rat-mult*: $\text{const-of-rat } (a * b) = \text{const-of-rat } a * \text{const-of-rat } b$
 $\langle \text{proof} \rangle$

lemma *const-of-rat-p-equal-0-iff*: $p\text{-equal } p (\text{const-of-rat } a) \ 0 \longleftrightarrow a = 0$
 $\langle \text{proof} \rangle$

end

lemma *const-of-rat-eq-0-iff* [simp]: $\text{const-of-rat } a = 0 \longleftrightarrow a = 0$
<proof>

lemma *const-of-rat-neg-one* [simp]: $\text{const-of-rat } (-1) = -1$
<proof>

lemma *const-of-rat-diff*: $\text{const-of-rat } (a - b) = \text{const-of-rat } a - \text{const-of-rat } b$
<proof>

lemma *const-of-rat-sum*: $\text{const-of-rat } (\sum a \in A. f a) = (\sum a \in A. \text{const-of-rat } (f a))$
<proof>

lemma *const-of-rat-prod*: $\text{const-of-rat } (\prod a \in A. f a) = (\prod a \in A. \text{const-of-rat } (f a))$
<proof>

lemma *const-of-rat-inverse*: $\text{const-of-rat } (\text{inverse } a) = \text{inverse } (\text{const-of-rat } a)$
<proof>

lemma *left-inverse-const-of-rat*[simp]:
 $\text{inverse } (\text{const-of-rat } a) * \text{const-of-rat } a = 1$ **if** $a \neq 0$
<proof>

lemma *right-inverse-const-of-rat*[simp]:
 $\text{const-of-rat } a * \text{inverse } (\text{const-of-rat } a) = 1$ **if** $a \neq 0$
<proof>

lemma *const-of-rat-divide*: $\text{const-of-rat } (a / b) = \text{const-of-rat } a / \text{const-of-rat } b$
<proof>

lemma *divide-self-const-of-rat*[simp]: $(\text{const-of-rat } a) / (\text{const-of-rat } a) = 1$ **if** $a \neq 0$
<proof>

lemma *const-of-rat-power*: $\text{const-of-rat } (a ^ n) = (\text{const-of-rat } a) ^ n$
<proof>

lemma *const-of-rat-eq-iff* [simp]: $\text{const-of-rat } a = \text{const-of-rat } b \longleftrightarrow a = b$
<proof>

lemma *zero-eq-const-of-rat-iff* [simp]: $0 = \text{const-of-rat } a \longleftrightarrow 0 = a$
<proof>

lemma *const-of-rat-eq-1-iff* [simp]: $\text{const-of-rat } a = 1 \longleftrightarrow a = 1$
<proof>

lemma *one-eq-const-of-rat-iff* [simp]: $1 = \text{const-of-rat } a \longleftrightarrow 1 = a$
<proof>

Collapse nested embeddings.

lemma *const-of-rat-of-nat-eq* [simp]: *const-of-rat (of-nat n) = of-nat n*
 ⟨proof⟩

lemma *const-of-rat-of-int-eq* [simp]: *const-of-rat (of-int z) = of-int z*
 ⟨proof⟩

Product of all p-adic embeddings of the field of rationals.

definition *range-const-of-rat* :: 'b set ($\mathbb{Q}_{\forall p}$)
 where $\mathbb{Q}_{\forall p} = \text{range const-of-rat}$

lemma *range-const-of-rat-cases* [cases set: *range-const-of-rat*]:
 assumes $q \in \mathbb{Q}_{\forall p}$
 obtains (*const-of-rat*) r where $q = \text{const-of-rat } r$
 ⟨proof⟩

lemma *range-const-of-rat-cases'*:
 assumes $x \in \mathbb{Q}_{\forall p}$
 obtains $a \ b$ where $b > 0$ coprime $a \ b$ $x = \text{of-int } a / \text{of-int } b$
 ⟨proof⟩

lemma *range-const-of-rat-of-rat* [simp]: *const-of-rat $r \in \mathbb{Q}_{\forall p}$*
 ⟨proof⟩

lemma *range-const-of-rat-of-int* [simp]: *of-int $z \in \mathbb{Q}_{\forall p}$*
 ⟨proof⟩

lemma *Ints-subset-range-const-of-rat*: $\mathbb{Z} \subseteq \mathbb{Q}_{\forall p}$
 ⟨proof⟩

lemma *range-const-of-rat-of-nat* [simp]: *of-nat $n \in \mathbb{Q}_{\forall p}$*
 ⟨proof⟩

lemma *Nats-subset-range-const-of-rat*: $\mathbb{N} \subseteq \mathbb{Q}_{\forall p}$
 ⟨proof⟩

lemma *range-const-of-rat-0* [simp]: *0 $\in \mathbb{Q}_{\forall p}$*
 ⟨proof⟩

lemma *range-const-of-rat-1* [simp]: *1 $\in \mathbb{Q}_{\forall p}$*
 ⟨proof⟩

lemma *range-const-of-rat-add* [simp]:
 $a \in \mathbb{Q}_{\forall p} \implies$
 $b \in \mathbb{Q}_{\forall p} \implies$
 $a + b \in \mathbb{Q}_{\forall p}$
 ⟨proof⟩

lemma *range-const-of-rat-minus-iff* [simp]:

$- a \in \mathbb{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathbb{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-rat-diff* [simp]:

$a \in \mathbb{Q}_{\forall p} \implies$
 $b \in \mathbb{Q}_{\forall p} \implies$
 $a - b \in \mathbb{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-rat-mult* [simp]:

$a \in \mathbb{Q}_{\forall p} \implies$
 $b \in \mathbb{Q}_{\forall p} \implies$
 $a * b \in \mathbb{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-rat-inverse* [simp]:

$a \in \mathbb{Q}_{\forall p} \implies$
 $inverse\ a \in \mathbb{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-rat-divide* [simp]:

$a \in \mathbb{Q}_{\forall p} \implies$
 $b \in \mathbb{Q}_{\forall p} \implies$
 $a / b \in \mathbb{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-rat-power* [simp]:

$a \in \mathbb{Q}_{\forall p} \implies$
 $a ^ n \in \mathbb{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-rat-sum* [intro]:

$(\bigwedge x. x \in A \implies f\ x \in \mathbb{Q}_{\forall p}) \implies$
 $sum\ f\ A \in \mathbb{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-rat-prod* [intro]:

$(\bigwedge x. x \in A \implies f\ x \in \mathbb{Q}_{\forall p})$
 $\implies prod\ f\ A \in \mathbb{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-rat-induct* [case-names const-of-rat, induct set: range-const-of-rat]:

$q \in \mathbb{Q}_{\forall p} \implies$
 $(\bigwedge r. P\ (const-of-rat\ r)) \implies P\ q$
 $\langle proof \rangle$

lemma *p-adic-Rats-p-equal-0-iff*:

$a \in \mathbb{Q}_{\forall p} \implies$

$p\text{-equal } p \ a \ 0 \longleftrightarrow a = 0$
 $\langle \text{proof} \rangle$

lemma *range-const-of-rat-infinite*: $\neg \text{finite } \mathbb{Q}_{\forall p}$
 $\langle \text{proof} \rangle$

lemma *left-inverse-range-const-of-rat*:
 $a \in \mathbb{Q}_{\forall p} \implies a \neq 0 \implies$
 $\text{inverse } a * a = 1$
 $\langle \text{proof} \rangle$

lemma *right-inverse-range-const-of-rat*:
 $a \in \mathbb{Q}_{\forall p} \implies a \neq 0 \implies$
 $a * \text{inverse } a = 1$
 $\langle \text{proof} \rangle$

lemma *divide-self-range-const-of-rat*:
 $a \in \mathbb{Q}_{\forall p} \implies a \neq 0 \implies$
 $a / a = 1$
 $\langle \text{proof} \rangle$

lemma *range-const-of-rat-add-iff*:
 $a \in \mathbb{Q}_{\forall p} \vee$
 $b \in \mathbb{Q}_{\forall p} \implies$
 $a + b \in \mathbb{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathbb{Q}_{\forall p} \wedge b \in \mathbb{Q}_{\forall p}$
 $\langle \text{proof} \rangle$

lemma *range-const-of-rat-diff-iff*:
 $a \in \mathbb{Q}_{\forall p} \vee$
 $b \in \mathbb{Q}_{\forall p} \implies$
 $a - b \in \mathbb{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathbb{Q}_{\forall p} \wedge b \in \mathbb{Q}_{\forall p}$
 $\langle \text{proof} \rangle$

lemma *range-const-of-rat-mult-iff*:
 $a * b \in \mathbb{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathbb{Q}_{\forall p} \wedge b \in \mathbb{Q}_{\forall p}$
if
 $a \in \mathbb{Q}_{\forall p} - \{0\} \vee$
 $b \in \mathbb{Q}_{\forall p} - \{0\}$
 $\langle \text{proof} \rangle$

lemma *range-const-of-rat-inverse-iff* [simp]:
 $\text{inverse } a \in \mathbb{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathbb{Q}_{\forall p}$
 $\langle \text{proof} \rangle$

lemma *range-const-of-rat-divide-iff*:

$a / b \in \mathbb{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathbb{Q}_{\forall p} \wedge b \in \mathbb{Q}_{\forall p}$
if
 $a \in \mathbb{Q}_{\forall p} - \{0\} \vee$
 $b \in \mathbb{Q}_{\forall p} - \{0\}$
 $\langle proof \rangle$

end

lemma *Fract-eq0-iff*: $Fraction-Field.Fract\ a\ b = 0 \longleftrightarrow a = 0 \vee b = 0$
 $\langle proof \rangle$

locale *global-p-equal-div-inv-w-inj-idom-consts* =
 $global-p-equal-div-inv-w-inj-consts\ p-equal\ p-restrict\ func-of-consts$
 $+ global-p-equality-w-inj-consts\ p-equal\ p-restrict\ func-of-consts$
for $p-equal$::
 $'a \Rightarrow 'b :: \{comm-ring-1, inverse, divide-trivial\} \Rightarrow$
 $'b \Rightarrow bool$
and $p-restrict$:: $'b \Rightarrow ('a \Rightarrow bool) \Rightarrow 'b$
and $func-of-consts$:: $('a \Rightarrow 'c :: idom) \Rightarrow 'b$
begin

lift-definition *const-of-fract* :: $'c\ fract \Rightarrow 'b$
is $\lambda x :: 'c \times 'c. const\ (fst\ x) / const\ (snd\ x)$
 $\langle proof \rangle$

lemma *inj-const-of-fract*: $inj\ const-of-fract$
 $\langle proof \rangle$

lemma *const-of-fract-0 [simp]*: $const-of-fract\ 0 = 0$
 $\langle proof \rangle$

lemma *const-of-fract-1 [simp]*: $const-of-fract\ 1 = 1$
 $\langle proof \rangle$

lemma *const-of-fract-Fract*: $const-of-fract\ (Fraction-Field.Fract\ a\ b) = const\ a /$
 $const\ b$
 $\langle proof \rangle$

lemma *const-of-fract-p-eq-0-iff [simp]*: $p-equal\ p\ (const-of-fract\ x)\ 0 \longleftrightarrow x = 0$
 $\langle proof \rangle$

lemma *const-of-fract-equal-0-iff [simp]*: $const-of-fract\ x = 0 \longleftrightarrow x = 0$
 $\langle proof \rangle$

lemma *const-of-fract-minus*: $const-of-fract\ (-\ a) = -\ const-of-fract\ a$
 $\langle proof \rangle$

lemma *const-of-fract-add*: $const-of-fract\ (a + b) = const-of-fract\ a + const-of-fract\ b$

b
 $\langle \text{proof} \rangle$

lemma *const-of-fract-mult*: $\text{const-of-fract } (a * b) = \text{const-of-fract } a * \text{const-of-fract } b$
 $\langle \text{proof} \rangle$

lemma *const-of-fract-neg-one* [simp]: $\text{const-of-fract } (- 1) = -1$
 $\langle \text{proof} \rangle$

lemma *const-of-fract-diff*: $\text{const-of-fract } (a - b) = \text{const-of-fract } a - \text{const-of-fract } b$
 $\langle \text{proof} \rangle$

lemma *const-of-fract-sum*:
 $\text{const-of-fract } (\sum a \in A. f a) = (\sum a \in A. \text{const-of-fract } (f a))$
 $\langle \text{proof} \rangle$

lemma *const-of-fract-prod*:
 $\text{const-of-fract } (\prod a \in A. f a) = (\prod a \in A. \text{const-of-fract } (f a))$
 $\langle \text{proof} \rangle$

lemma *const-of-fract-inverse*: $\text{const-of-fract } (\text{inverse } a) = \text{inverse } (\text{const-of-fract } a)$
 $\langle \text{proof} \rangle$

lemma *left-inverse-const-of-fract*[simp]:
 $\text{inverse } (\text{const-of-fract } a) * \text{const-of-fract } a = 1 \text{ if } a \neq 0$
 $\langle \text{proof} \rangle$

lemma *right-inverse-const-of-fract*[simp]:
 $(\text{const-of-fract } a) * \text{inverse } (\text{const-of-fract } a) = 1 \text{ if } a \neq 0$
 $\langle \text{proof} \rangle$

lemma *const-of-fract-divide*: $\text{const-of-fract } (a / b) = \text{const-of-fract } a / \text{const-of-fract } b$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-divide-self-of-fract*[simp]:
 $(\text{const-of-fract } a) / (\text{const-of-fract } a) = 1 \text{ if } a \neq 0$
 $\langle \text{proof} \rangle$

lemma *const-of-fract-power*: $\text{const-of-fract } (a ^ n) = (\text{const-of-fract } a) ^ n$
 $\langle \text{proof} \rangle$

lemma *const-of-fract-eq-iff* [simp]:
 $\text{const-of-fract } a = \text{const-of-fract } b \longleftrightarrow a = b$
 $\langle \text{proof} \rangle$

lemma *zero-eq-const-of-fract-iff* [simp]: $0 = \text{const-of-fract } a \longleftrightarrow 0 = a$
 ⟨proof⟩

lemma *const-of-fract-eq-1-iff* [simp]: $\text{const-of-fract } a = 1 \longleftrightarrow a = 1$
 ⟨proof⟩

lemma *one-eq-const-of-fract-iff* [simp]: $1 = \text{const-of-fract } a \longleftrightarrow 1 = a$
 ⟨proof⟩

Collapse nested embeddings.

lemma *const-of-fract-numer-eq-const* [simp]: $\text{const-of-fract } (\text{Fraction-Field.Fract } a \ 1) = \text{const } a$
 ⟨proof⟩

lemma *const-of-fract-of-nat-eq* [simp]: $\text{const-of-fract } (\text{of-nat } n) = \text{of-nat } n$
 ⟨proof⟩

lemma *const-of-fract-of-int-eq* [simp]: $\text{const-of-fract } (\text{of-int } z) = \text{of-int } z$
 ⟨proof⟩

Product of all p-adic embeddings of the field of fractions of the base ring.

definition *range-const-of-fract* :: 'b set ($\mathbb{Q}_{\forall p}$)
 where $\mathbb{Q}_{\forall p} = \text{range const-of-fract}$

lemma *range-const-of-fract-cases* [cases set: *range-const-of-fract*]:
 assumes $q \in \mathbb{Q}_{\forall p}$
 obtains $(\text{const-of-fract}) \ r$ where $q = \text{const-of-fract } r$
 ⟨proof⟩

lemma *range-const-of-fract-cases'*:
 assumes $x \in \mathbb{Q}_{\forall p}$
 obtains $a \ b$ where $b \neq 0 \ x = \text{const } a / \text{const } b$
 ⟨proof⟩

lemma *range-const-of-fract-of-fract* [simp]:
 $\text{const-of-fract } r \in \mathbb{Q}_{\forall p}$
 ⟨proof⟩

lemma *range-const-of-fract-const* [simp]: $\text{const } a \in \mathbb{Q}_{\forall p}$
 ⟨proof⟩

lemma *range-const-of-fract-of-int* [simp]: $\text{of-int } z \in \mathbb{Q}_{\forall p}$
 ⟨proof⟩

lemma *Ints-subset-range-const-of-fract*: $\mathbb{Z} \subseteq \mathbb{Q}_{\forall p}$
 ⟨proof⟩

lemma *range-const-of-fract-of-nat* [simp]: $\text{of-nat } n \in \mathbb{Q}_{\forall p}$
 ⟨proof⟩

lemma *Nats-subset-range-const-of-fract*: $\mathbb{N} \subseteq \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-0* [simp]: $0 \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-1* [simp]: $1 \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-add* [simp]:
 $a \in \mathcal{Q}_{\forall p} \implies$
 $b \in \mathcal{Q}_{\forall p} \implies$
 $a + b \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-minus-iff* [simp]:
 $- a \in \mathcal{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-diff* [simp]:
 $a \in \mathcal{Q}_{\forall p} \implies$
 $b \in \mathcal{Q}_{\forall p} \implies$
 $a - b \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-mult* [simp]:
 $a \in \mathcal{Q}_{\forall p} \implies$
 $b \in \mathcal{Q}_{\forall p} \implies$
 $a * b \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-inverse* [simp]:
 $a \in \mathcal{Q}_{\forall p} \implies$
 $inverse\ a \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-divide* [simp]:
 $a \in \mathcal{Q}_{\forall p} \implies$
 $b \in \mathcal{Q}_{\forall p} \implies$
 $a / b \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-power* [simp]:
 $a \in \mathcal{Q}_{\forall p} \implies$
 $a ^ n \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-sum* [intro]:

$(\bigwedge x. x \in A \implies$
 $f\ x \in \mathcal{Q}_{\forall p}) \implies$
 $sum\ f\ A \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract* [intro]:

$(\bigwedge x. x \in A \implies f\ x \in \mathcal{Q}_{\forall p})$
 $\implies prod\ f\ A \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-induct* [case-names const-of-fract, induct set: range-const-of-fract]:

$q \in \mathcal{Q}_{\forall p} \implies$
 $(\bigwedge r. P\ (const-of-fract\ r)) \implies P\ q$
 $\langle proof \rangle$

lemma *p-adic-Fracts-p-equal-0-iff*:

$a \in \mathcal{Q}_{\forall p} \implies$
 $p\text{-equal}\ p\ a\ 0 \longleftrightarrow a = 0$
 $\langle proof \rangle$

lemma *range-const-of-fract-infinite*:

infinite $\mathcal{Q}_{\forall p}$ **if** *infinite* (*UNIV* :: 'c set)
 $\langle proof \rangle$

lemma *left-inverse-range-const-of-fract*:

$a \in \mathcal{Q}_{\forall p} \implies a \neq 0 \implies$
 $inverse\ a * a = 1$
 $\langle proof \rangle$

lemma *right-inverse-range-const-of-fract*:

$a \in \mathcal{Q}_{\forall p} \implies a \neq 0 \implies$
 $a * inverse\ a = 1$
 $\langle proof \rangle$

lemma *divide-self-range-const-of-fract*:

$a \in \mathcal{Q}_{\forall p} \implies a \neq 0 \implies$
 $a / a = 1$
 $\langle proof \rangle$

lemma *range-const-of-fract-add-iff*:

$a \in \mathcal{Q}_{\forall p} \vee$
 $b \in \mathcal{Q}_{\forall p} \implies$
 $a + b \in \mathcal{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathcal{Q}_{\forall p} \wedge b \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-diff-iff*:

$a \in \mathcal{Q}_{\forall p} \vee$

$b \in \mathcal{Q}_{\forall p} \implies$
 $a - b \in \mathcal{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathcal{Q}_{\forall p} \wedge b \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-mult-iff*:

$a * b \in \mathcal{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathcal{Q}_{\forall p} \wedge b \in \mathcal{Q}_{\forall p}$
if
 $a \in \mathcal{Q}_{\forall p} - \{0\} \vee$
 $b \in \mathcal{Q}_{\forall p} - \{0\}$
 $\langle proof \rangle$

lemma *range-const-of-fract-inverse-iff* [simp]:

$inverse\ a \in \mathcal{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathcal{Q}_{\forall p}$
 $\langle proof \rangle$

lemma *range-const-of-fract-divide-iff*:

$a / b \in \mathcal{Q}_{\forall p} \longleftrightarrow$
 $a \in \mathcal{Q}_{\forall p} \wedge b \in \mathcal{Q}_{\forall p}$
if
 $a \in \mathcal{Q}_{\forall p} - \{0\} \vee$
 $b \in \mathcal{Q}_{\forall p} - \{0\}$
 $\langle proof \rangle$

end

locale *global-p-equality-w-inj-index-consts* = *global-p-equality-w-inj-consts*

+

fixes *index-inj* :: 'a $\Rightarrow$ 'c
assumes *index-inj* : *inj index-inj*
and *index-inj-nzero*: *index-inj p $\neq$ 0*

begin

abbreviation *global-uniformizer* $\equiv$ *consts index-inj*

abbreviation *index-const p* $\equiv$ *const (index-inj p)*

lemma *global-uniformizer-locally-uniformizes*: *p-equal p global-uniformizer (index-const p)*

$\langle proof \rangle$

lemma *index-const-globally-nequiv-0*: $\neg$ *p-equal q (index-const p) 0*

$\langle proof \rangle$

lemma *global-uniformizer-globally-nequiv-0*: $\neg$ *p-equal p global-uniformizer 0*

$\langle proof \rangle$

end

```

locale global-p-equality-no-zero-divisors-1-w-inj-index-consts =
  global-p-equality-w-inj-index-consts + p-equality-no-zero-divisors-1
begin

lemma index-const-pow-globally-nequiv-0:  $\neg p\text{-equal } q \text{ (index-const } p \wedge n) \ 0$ 
  <proof>

lemma global-uniformizer-pow-globally-nequiv-0:  $\neg p\text{-equal } q \text{ (global-uniformizer } \wedge n) \ 0$ 
  <proof>

end

locale global-p-equality-depth-w-inj-index-consts =
  p-equality-depth p-equal p-depth
+ global-p-equality-w-inj-index-consts p-equal p-restrict func-of-consts index-inj
  for p-equal      :: 'a  $\Rightarrow$  'b::comm-ring-1  $\Rightarrow$  'b  $\Rightarrow$  bool
  and p-restrict   :: 'b  $\Rightarrow$  ('a  $\Rightarrow$  bool)  $\Rightarrow$  'b
  and p-depth      :: 'a  $\Rightarrow$  'b  $\Rightarrow$  int
  and func-of-consts :: ('a  $\Rightarrow$  'c::{factorial-semiring,ring-1})  $\Rightarrow$  'b
  and index-inj    :: 'a  $\Rightarrow$  'c
+
  assumes p-depth-func-of-consts:
    p-depth p (func-of-consts f) = int (multiplicity (index-inj p) (f p))
  and index-inj-nunit :  $\neg$  is-unit (index-inj p)
  and index-inj-coprime:
    p  $\neq$  q  $\implies$  multiplicity (index-inj p) (index-inj q) = 0
begin

lemma p-depth-set-consts: func-of-consts f  $\in$  p-depth-set p 0
  <proof>

lemma global-depth-set-consts: func-of-consts f  $\in$  global-depth-set 0
  <proof>

lemma p-depth-1[simp]: p-depth p (1::'b) = 0
  <proof>

lemma p-depth-index-const: p-depth p (index-const p) = 1
  <proof>

lemma p-depth-index-const': p-depth q (index-const p) = of_bool (q = p)
  <proof>

lemma p-depth-global-uniformizer: p-depth p (global-uniformizer::'b) = 1
  <proof>

lemma p-depth-consts-func-ge0: p-depth p (consts f)  $\geq$  0

```

```

    <proof>

end

locale global-p-equal-depth-no-zero-divisors-w-inj-index-consts =
  global-p-equal p-equal p-restrict
+ global-p-equality-depth-no-zero-divisors-1 p-equal p-restrict p-depth
+ global-p-equality-depth-w-inj-index-consts p-equal p-restrict p-depth func-of-consts
index-inj
  for p-equal      :: 'a  $\Rightarrow$  'b::comm-ring-1  $\Rightarrow$  'b  $\Rightarrow$  bool
  and p-restrict   :: 'b  $\Rightarrow$  ('a  $\Rightarrow$  bool)  $\Rightarrow$  'b
  and p-depth      :: 'a  $\Rightarrow$  'b  $\Rightarrow$  int
  and func-of-consts :: ('a  $\Rightarrow$  'c::{factorial-semiring,ring-1})  $\Rightarrow$  'b
  and index-inj    :: 'a  $\Rightarrow$  'c
begin

sublocale global-p-equal-w-inj-consts <proof>
sublocale global-p-equal-depth-no-zero-divisors <proof>

lemma p-depth-index-const-pow: p-depth p ((index-const p) ^ n) = int n
  <proof>

lemma p-depth-index-const-pow': p-depth q ((index-const p) ^ n) = int n * of-bool
(q = p)
  <proof>

lemma p-depth-global-uniformizer-pow: p-depth p (global-uniformizer ^ n) = int n
  <proof>

end

locale global-p-equal-depth-div-inv-w-inj-index-consts =
  global-p-equal-depth-div-inv p-equal p-depth p-restrict
+ global-p-equal-depth-no-zero-divisors-w-inj-index-consts
  p-equal p-restrict p-depth func-of-consts index-inj
  for p-equal      ::
    'a  $\Rightarrow$  'b::{comm-ring-1, inverse, divide-trivial}  $\Rightarrow$ 
    'b  $\Rightarrow$  bool
  and p-depth      :: 'a  $\Rightarrow$  'b  $\Rightarrow$  int
  and p-restrict   :: 'b  $\Rightarrow$  ('a  $\Rightarrow$  bool)  $\Rightarrow$  'b
  and func-of-consts :: ('a  $\Rightarrow$  'c::{factorial-semiring,ring-1})  $\Rightarrow$  'b
  and index-inj    :: 'a  $\Rightarrow$  'c
begin

lemma index-const-powi-globally-nequiv-0:  $\neg$  p-equal q (index-const p powi n) 0
  <proof>

lemma global-uniformizer-powi-globally-nequiv-0:  $\neg$  p-equal q (global-uniformizer

```

$\text{powi } n) \ 0$
 $\langle \text{proof} \rangle$

lemma *index-const-powi-add*:
 $(\text{index-const } p) \ \text{powi } (m + n) = (\text{index-const } p) \ \text{powi } m * (\text{index-const } p) \ \text{powi } n$
 $\langle \text{proof} \rangle$

lemma *global-uniformizer-powi-add*:
 $\text{global-uniformizer } \text{powi } (m + n) = \text{global-uniformizer } \text{powi } m * \text{global-uniformizer } \text{powi } n$
 $\langle \text{proof} \rangle$

lemma *p-depth-index-const-powi*: $p\text{-depth } p \ (\text{index-const } p \ \text{powi } n) = n$
 $\langle \text{proof} \rangle$

lemma *p-depth-index-const-powi'*: $p\text{-depth } q \ (\text{index-const } p \ \text{powi } n) = n * \text{of-bool } (q = p)$
 $\langle \text{proof} \rangle$

lemma *p-depth-global-uniformizer-powi*: $p\text{-depth } p \ (\text{global-uniformizer } \text{powi } n) = n$
 $\langle \text{proof} \rangle$

lemma *local-uniformizer-locally-uniformizes*:
 $p\text{-depth } p \ ((\text{index-const } p) \ \text{powi } (-n) * x) = 0 \ \text{if } p\text{-depth } p \ x = n$
 $\langle \text{proof} \rangle$

lemma *global-uniformizer-locally-uniformizes*:
 $p\text{-depth } p \ (\text{global-uniformizer } \text{powi } (-n) * x) = 0 \ \text{if } p\text{-depth } p \ x = n$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-eq-index-const-coset*:
 $p\text{-depth-set } p \ n = ((\text{index-const } p) \ \text{powi } n) * o \ (p\text{-depth-set } p \ 0)$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-eq-index-const-coset'*:
 $p\text{-depth-set } p \ 0 = ((\text{index-const } p) \ \text{powi } (-n)) * o \ (p\text{-depth-set } p \ n)$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-eq-global-uniformizer-coset*:
 $p\text{-depth-set } p \ n = (\text{global-uniformizer } \text{powi } n) * o \ (p\text{-depth-set } p \ 0)$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-eq-global-uniformizer-coset'*:
 $p\text{-depth-set } p \ 0 = (\text{global-uniformizer } \text{powi } (-n)) * o \ (p\text{-depth-set } p \ n)$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-eq-global-uniformizer-coset*:
 $\text{global-depth-set } n = (\text{global-uniformizer } \text{powi } n) * o \ (\text{global-depth-set } 0)$

$\langle \text{proof} \rangle$

lemma *global-depth-set-eq-global-uniformizer-coset'*:

global-depth-set 0 = (*global-uniformizer* powi (-n)) *o (*global-depth-set* n)

$\langle \text{proof} \rangle$

definition *shift-p-depth* :: 'a $\Rightarrow$ int $\Rightarrow$ 'b $\Rightarrow$ 'b

where *shift-p-depth* p n x = x * (consts (λq . if q = p then index-inj p else 1))
powi n

lemma *shift-p-depth-0[simp]*: *shift-p-depth* p n 0 = 0

$\langle \text{proof} \rangle$

lemma *shift-p-depth-by-0[simp]*: *shift-p-depth* p 0 x = x

$\langle \text{proof} \rangle$

lemma *shift-p-depth-add[simp]*:

shift-p-depth p n (x + y) = *shift-p-depth* p n x + *shift-p-depth* p n y

$\langle \text{proof} \rangle$

lemma *shift-p-depth-uminus[simp]*:

shift-p-depth p n (- x) = - *shift-p-depth* p n x

$\langle \text{proof} \rangle$

lemma *shift-p-depth-minus[simp]*:

shift-p-depth p n (x - y) = *shift-p-depth* p n x - *shift-p-depth* p n y

$\langle \text{proof} \rangle$

lemma *shift-p-depth-equiv-at-place*: p-equal p (*shift-p-depth* p n x) (x * (index-const p powi n))

$\langle \text{proof} \rangle$

lemma *shift-p-depth-equiv-at-other-places*:

p-equal q (*shift-p-depth* p n x) x **if** q $\neq$ p

$\langle \text{proof} \rangle$

lemma *shift-p-depth-equiv0-iff*:

p-equal q (*shift-p-depth* p n x) 0 $\longleftrightarrow$ p-equal q x 0

$\langle \text{proof} \rangle$

lemma *shift-p-depth-equiv-iff*:

p-equal q (*shift-p-depth* p n x) (*shift-p-depth* p n y) = p-equal q x y

$\langle \text{proof} \rangle$

lemma *shift-p-depth-trivial-iff*:

shift-p-depth p n x = x $\longleftrightarrow$ n = 0 $\vee$ p-equal p x 0

$\langle \text{proof} \rangle$

lemma *shift-p-depth-times-right*: *shift-p-depth* p n (x * y) = x * *shift-p-depth* p n

y
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-times-left*: $\text{shift-p-depth } p \ n \ (x * y) = \text{shift-p-depth } p \ n \ x * y$
 $\langle \text{proof} \rangle$

lemma *shift-shift-p-depth*: $\text{shift-p-depth } p \ n \ (\text{shift-p-depth } p \ m \ x) = \text{shift-p-depth } p \ (m + n) \ x$
 $\langle \text{proof} \rangle$

lemma *shift-shift-p-depth-image*:
 $\text{shift-p-depth } p \ n \ ' \ \text{shift-p-depth } p \ m \ ' \ A = \text{shift-p-depth } p \ (m + n) \ ' \ A$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-at-place*:
 $p\text{-depth } p \ (\text{shift-p-depth } p \ n \ x) = p\text{-depth } p \ x + n \ \text{if } n = 0 \vee \neg p\text{-equal } p \ x \ 0$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-at-other-places*:
 $p\text{-depth } q \ (\text{shift-p-depth } p \ n \ x) = p\text{-depth } q \ x \ \text{if } q \neq p$
 $\langle \text{proof} \rangle$

lemma *p-depth-set-shift-p-depth-closed*:
 $\text{shift-p-depth } p \ n \ x \in p\text{-depth-set } p \ m \ \text{if } p\text{-depth } p \ x \geq m - n$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-shift-p-depth-closed*:
 $\text{shift-p-depth } p \ n \ x \in \text{global-depth-set } m$
if $x \in \text{global-depth-set } m$
and $\neg p\text{-equal } p \ x \ 0 \longrightarrow p\text{-depth } p \ x \geq m - n$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-mem-p-depth-set*:
 $p\text{-depth } p \ x \geq m - n$
if $\neg p\text{-equal } p \ x \ 0$ **and** $\text{shift-p-depth } p \ n \ x \in p\text{-depth-set } p \ m$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-mem-global-depth-set*:
 $p\text{-depth } p \ x \geq m - n$
if $\neg p\text{-equal } p \ x \ 0$ **and** $\text{shift-p-depth } p \ n \ x \in \text{global-depth-set } m$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-p-depth-set*: $\text{shift-p-depth } p \ n \ ' \ p\text{-depth-set } p \ m = p\text{-depth-set } p \ (m + n)$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-cong*:
 $p\text{-equal } q \ (\text{shift-p-depth } p \ n \ x) \ (\text{shift-p-depth } p \ n \ y) \ \text{if } p\text{-equal } q \ x \ y$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-p-restrict*:

shift-p-depth p n (p-restrict x P) = p-restrict (shift-p-depth p n x) P
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-p-restrict-global-depth-set-memI*:

shift-p-depth p n (p-restrict x ((=) p)) $\in$ global-depth-set m
if $\neg p\text{-equal } p \ x \ 0 \longrightarrow p\text{-depth } p \ x \geq m - n$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-p-restrict-global-depth-set-memI'*:

shift-p-depth p n (p-restrict x ((=) p)) $\in$ global-depth-set (m + n)
if $x \in p\text{-depth-set } p \ m$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-p-restrict-global-depth-set-image*:

shift-p-depth p n ' ($\lambda x. p\text{-restrict } x ((=) p)$) ' global-depth-set m
 $= (\lambda x. p\text{-restrict } x ((=) p)) ' \text{global-depth-set } (m + n)$
 $\langle \text{proof} \rangle$

end

3.7 Topological patterns

3.7.1 By place

Convergence of sequences as measured by depth

context *p-equality-depth*
begin

definition *p-limseq-condition* ::

$'a \Rightarrow (\text{nat} \Rightarrow 'b) \Rightarrow 'b \Rightarrow \text{int} \Rightarrow$
 $\text{nat} \Rightarrow \text{bool}$

where

p-limseq-condition p X x n K =
 $(\forall k \geq K. \neg p\text{-equal } p \ (X \ k) \ x \longrightarrow p\text{-depth } p \ (X \ k - x) > n)$

lemma *p-limseq-conditionI* [intro]:

p-limseq-condition p X x n K
if
 $\bigwedge k. k \geq K \implies \neg p\text{-equal } p \ (X \ k) \ x \implies$
 $p\text{-depth } p \ (X \ k - x) > n$
 $\langle \text{proof} \rangle$

lemma *p-limseq-conditionD*:

p-depth p (X k - x) > n
if *p-limseq-condition p X x n K* **and** $k \geq K$ **and** $\neg p\text{-equal } p \ (X \ k) \ x$
 $\langle \text{proof} \rangle$

abbreviation $p\text{-limseq} ::$

$'a \Rightarrow (\text{nat} \Rightarrow 'b) \Rightarrow 'b \Rightarrow \text{bool}$

where $p\text{-limseq } p \ X \ x \equiv (\forall n. \exists K. p\text{-limseq-condition } p \ X \ x \ n \ K)$

lemma $p\text{-limseq-p-cong}$:

$p\text{-limseq } p \ X \ x$

if $\forall n \geq N. p\text{-equal } p \ (X \ n) \ (Y \ n)$ **and** $p\text{-equal } p \ x \ y$ **and** $p\text{-limseq } p \ Y \ y$
 $\langle \text{proof} \rangle$

lemma $p\text{-limseq-p-constant}$: $p\text{-limseq } p \ X \ x$ **if** $\forall n. p\text{-equal } p \ (X \ n) \ x$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-constant}$: $p\text{-limseq } p \ (\lambda n. x) \ x$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-p-eventually-constant}$:

$p\text{-limseq } p \ X \ x$ **if** $\forall_F k \text{ in sequentially. } p\text{-equal } p \ (X \ k) \ x$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-0 [simp]}$: $p\text{-limseq } p \ 0 \ 0$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-0-iff}$: $p\text{-limseq } p \ 0 \ x \longleftrightarrow p\text{-equal } p \ x \ 0$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-add}$: $p\text{-limseq } p \ (X + Y) \ (x + y)$ **if** $p\text{-limseq } p \ X \ x$ **and** $p\text{-limseq}$

$p \ Y \ y$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-uminus}$: $p\text{-limseq } p \ (-X) \ (-x)$ **if** $p\text{-limseq } p \ X \ x$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-diff}$: $p\text{-limseq } p \ (X - Y) \ (x - y)$ **if** $p\text{-limseq } p \ X \ x$ **and** $p\text{-limseq}$

$p \ Y \ y$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-conv-0}$: $p\text{-limseq } p \ X \ x \longleftrightarrow p\text{-limseq } p \ (\lambda n. X \ n - x) \ 0$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-unique}$: $p\text{-equal } p \ x \ y$ **if** $p\text{-limseq } p \ X \ x$ **and** $p\text{-limseq } p \ X \ y$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-eventually-nequiv-0}$:

$\forall_F k \text{ in sequentially. } \neg p\text{-equal } p \ (X \ k) \ 0$

if $\neg p\text{-equal } p \ x \ 0$ **and** $p\text{-limseq } p \ X \ x$

$\langle \text{proof} \rangle$

lemma $p\text{-limseq-bdd-depth}$:

assumes $\neg p\text{-equal } p \ x \ 0$ **and** $p\text{-limseq } p \ X \ x$

obtains d **where** $\forall k. \neg p\text{-equal } p (X k) 0 \longrightarrow p\text{-depth } p (X k) \leq d$
 $\langle \text{proof} \rangle$

lemma $p\text{-limseq-eventually-bdd-depth}$:
 $\forall_F k \text{ in sequentially. } p\text{-depth } p (X k) \leq p\text{-depth } p x$
if $\neg p\text{-equal } p x 0$ **and** $p\text{-limseq } p X x$
 $\langle \text{proof} \rangle$

lemma $p\text{-limseq-delayed-iff}$:
 $p\text{-limseq } p X x \longleftrightarrow p\text{-limseq } p (\lambda n. X (N + n)) x$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-set-}p\text{-limseq-closed}$:
 $x \in p\text{-depth-set } p n$
if $p\text{-limseq } p X x$ **and** $\forall_F k \text{ in sequentially. } X k \in p\text{-depth-set } p n$
 $\langle \text{proof} \rangle$

Cauchy sequences by place

definition $p\text{-cauchy-condition} ::$
 $'a \Rightarrow (\text{nat} \Rightarrow 'b) \Rightarrow \text{int} \Rightarrow \text{nat} \Rightarrow \text{bool}$
where
 $p\text{-cauchy-condition } p X n K =$
 $(\forall k1 \geq K. \forall k2 \geq K.$
 $\neg p\text{-equal } p (X k1) (X k2) \longrightarrow p\text{-depth } p (X k1 - X k2) > n$
 $)$

lemma $p\text{-cauchy-conditionI}$:
 $p\text{-cauchy-condition } p X n K$
if
 $\bigwedge k k'. k \geq K \implies k' \geq K \implies$
 $\neg p\text{-equal } p (X k) (X k') \implies p\text{-depth } p (X k - X k') > n$
 $\langle \text{proof} \rangle$

lemma $p\text{-cauchy-conditionD}$:
 $p\text{-depth } p (X k - X k') > n$
if $p\text{-cauchy-condition } p X n K$ **and** $k \geq K$ **and** $k' \geq K$
and $\neg p\text{-equal } p (X k) (X k')$
 $\langle \text{proof} \rangle$

lemma $p\text{-cauchy-condition-mono-seq-tail}$:
 $p\text{-cauchy-condition } p X n K \implies p\text{-cauchy-condition } p X n K'$
if $K' \geq K$
 $\langle \text{proof} \rangle$

lemma $p\text{-cauchy-condition-mono-depth}$:
 $p\text{-cauchy-condition } p X n K \implies p\text{-cauchy-condition } p X m K$
if $m \leq n$
 $\langle \text{proof} \rangle$

abbreviation $p\text{-cauchy } p \ X \equiv (\forall n. \exists K. p\text{-cauchy-condition } p \ X \ n \ K)$

lemma $p\text{-cauchy-condition-LEAST}$:

$p\text{-cauchy-condition } p \ X \ n \ (LEAST \ K. p\text{-cauchy-condition } p \ X \ n \ K)$

if $p\text{-cauchy } p \ X$

$\langle proof \rangle$

lemma $p\text{-cauchy-condition-LEAST-mono}$:

$(LEAST \ K. p\text{-cauchy-condition } p \ X \ m \ K) \leq (LEAST \ K. p\text{-cauchy-condition } p \ X \ n \ K)$

if $p\text{-cauchy } p \ X$ **and** $m \leq n$

$\langle proof \rangle$

definition $p\text{-consec-cauchy-condition} ::$

$'a \Rightarrow (nat \Rightarrow 'b) \Rightarrow int \Rightarrow nat \Rightarrow bool$

where

$p\text{-consec-cauchy-condition } p \ X \ n \ K =$

$(\forall k \geq K.$

$\neg p\text{-equal } p \ (X \ (Suc \ k)) \ (X \ k) \longrightarrow p\text{-depth } p \ (X \ (Suc \ k) - X \ k) > n$

$)$

lemma $p\text{-consec-cauchy-conditionI}$:

$p\text{-consec-cauchy-condition } p \ X \ n \ K$

if

$\bigwedge k. k \geq K \implies \neg p\text{-equal } p \ (X \ (Suc \ k)) \ (X \ k) \implies$

$p\text{-depth } p \ (X \ (Suc \ k) - X \ k) > n$

$\langle proof \rangle$

lemma $p\text{-consec-cauchy}$:

$p\text{-cauchy } p \ X = (\forall n. \exists K. p\text{-consec-cauchy-condition } p \ X \ n \ K)$

$\langle proof \rangle$

end

context $p\text{-equality-depth-no-zero-divisors}$

begin

lemma $p\text{-limseq-mult-both-0}$: $p\text{-limseq } p \ (X * Y) \ 0$ **if** $p\text{-limseq } p \ X \ 0$ **and** $p\text{-limseq } p \ Y \ 0$

$\langle proof \rangle$

lemma $p\text{-limseq-constant-mult-0}$: $p\text{-limseq } p \ (\lambda n. x * Y \ n) \ 0$ **if** $p\text{-limseq } p \ Y \ 0$

$\langle proof \rangle$

lemma $p\text{-limseq-mult-0-right}$: $p\text{-limseq } p \ (X * Y) \ 0$ **if** $p\text{-limseq } p \ X \ x$ **and** $p\text{-limseq } p \ Y \ 0$

$\langle proof \rangle$

lemma $p\text{-limseq-mult}$: $p\text{-limseq } p \ (X * Y) \ (x * y)$ **if** $p\text{-limseq } p \ X \ x$ **and** $p\text{-limseq } p \ Y \ y$

$p \ Y \ y$
 $\langle proof \rangle$

lemma *poly-continuous-p-open-nbhds*:
 $p\text{-limseq } p \ (\lambda n. \text{poly } f \ (X \ n)) \ (\text{poly } f \ x) \text{ if } p\text{-limseq } p \ X \ x$
 $\langle proof \rangle$

end

context *p-equality-depth-no-zero-divisors-1*
begin

lemma *p-limseq-pow*: $p\text{-limseq } p \ (\lambda k. (X \ k)^\wedge n) \ (x^\wedge n) \text{ if } p\text{-limseq } p \ X \ x$
 $\langle proof \rangle$

lemma *p-limseq-poly*: $p\text{-limseq } p \ (\lambda n. \text{poly } f \ (X \ n)) \ (\text{poly } f \ x) \text{ if } p\text{-limseq } p \ X \ x$
 $\langle proof \rangle$

end

context *p-equality-depth-div-inv*
begin

lemma *p-limseq-inverse*:
 $p\text{-limseq } p \ (\lambda n. \text{inverse } (X \ n)) \ (\text{inverse } x)$
if $\neg p\text{-equal } p \ x \ 0$ **and** $p\text{-limseq } p \ X \ x$
 $\langle proof \rangle$

lemma *p-limseq-divide*:
 $p\text{-limseq } p \ (\lambda n. X \ n \ / \ Y \ n) \ (x \ / \ y)$
if $\neg p\text{-equal } p \ y \ 0$ **and** $p\text{-limseq } p \ X \ x$ **and** $p\text{-limseq } p \ Y \ y$
 $\langle proof \rangle$

end

context *global-p-equality-depth*
begin

lemma *p-limseq-p-restrict-iff*:
 $p\text{-limseq } p \ (\lambda n. p\text{-restrict } (X \ n) \ P) \ x \longleftrightarrow p\text{-limseq } p \ X \ x \text{ if } P \ p$
 $\langle proof \rangle$

lemma *p-limseq-p-restricted*: $p\text{-limseq } p \ (\lambda n. p\text{-restrict } (X \ n) \ P) \ 0 \text{ if } \neg P \ p$
 $\langle proof \rangle$

end

Base of open sets at local places

context *global-p-equal-depth*

begin

abbreviation $p\text{-open-nbhd} :: 'a \Rightarrow \text{int} \Rightarrow 'b \Rightarrow 'b \text{ set}$
where $p\text{-open-nbhd } p \ n \ x \equiv x + o \ p\text{-depth-set } p \ n$

abbreviation $\text{local-}p\text{-open-nbhds } p \equiv \{ p\text{-open-nbhd } p \ n \ x \mid n \ x. \text{ True} \}$
abbreviation $p\text{-open-nbhds} \equiv (\bigcup p. \text{local-}p\text{-open-nbhds } p)$

lemma $p\text{-open-nbhd}$: $x \in p\text{-open-nbhd } p \ n \ x$
 $\langle \text{proof} \rangle$

lemma $\text{local-}p\text{-open-nbhd-is-open}$: $\text{generate-topology } (\text{local-}p\text{-open-nbhds } p) \ (p\text{-open-nbhd } p \ n \ x)$
 $\langle \text{proof} \rangle$

lemma $p\text{-open-nbhd-is-open}$: $\text{generate-topology } p\text{-open-nbhds} \ (p\text{-open-nbhd } p \ n \ x)$
 $\langle \text{proof} \rangle$

lemma $\text{local-}p\text{-open-nbhds-are-coarser}$:
 $\text{generate-topology } (\text{local-}p\text{-open-nbhds } p) \ S \implies \text{generate-topology } p\text{-open-nbhds } S$
 $\langle \text{proof} \rangle$

lemma $p\text{-open-nbhd-diff-depth}$:
 $p\text{-depth } p \ (y - x) \geq n \text{ if } \neg p\text{-equal } p \ y \ x \text{ and } y \in p\text{-open-nbhd } p \ n \ x$
 $\langle \text{proof} \rangle$

lemma $p\text{-open-nbhd-eq-circle}$:
 $p\text{-open-nbhd } p \ n \ x = \{y. \neg p\text{-equal } p \ y \ x \longrightarrow p\text{-depth } p \ (y - x) \geq n\}$
 $\langle \text{proof} \rangle$

lemma $p\text{-open-nbhd-circle-multicentre}$:
 $p\text{-open-nbhd } p \ n \ y = p\text{-open-nbhd } p \ n \ x \text{ if } y \in p\text{-open-nbhd } p \ n \ x$
 $\langle \text{proof} \rangle$

lemma $p\text{-open-nbhd-circle-multicentre}'$:
 $x \in p\text{-open-nbhd } p \ n \ y \longleftrightarrow y \in p\text{-open-nbhd } p \ n \ x$
 $\langle \text{proof} \rangle$

lemma $p\text{-open-nbhd-}p\text{-restrict}$:
 $p\text{-open-nbhd } p \ n \ (p\text{-restrict } x \ P) = p\text{-open-nbhd } p \ n \ x \text{ if } P \ p$
 $\langle \text{proof} \rangle$

lemma $p\text{-restrict-}p\text{-open-nbhd-mem-iff}$:
 $y \in p\text{-open-nbhd } p \ n \ x \longleftrightarrow p\text{-restrict } y \ P \in p\text{-open-nbhd } p \ n \ x \text{ if } P \ p$
 $\langle \text{proof} \rangle$

lemma $p\text{-open-nbhd-}p\text{-restrict-add-mixed-drop-right}$:
 $p\text{-open-nbhd } p \ n \ (p\text{-restrict } x \ P + p\text{-restrict } y \ Q) = p\text{-open-nbhd } p \ n \ x$
if $P \ p$ **and** $\neg Q \ p$

$\langle \text{proof} \rangle$

lemma *p-open-nbhd-p-restrict-add-mixed-drop-left*:

$p\text{-open-nbhd } p \ n \ (p\text{-restrict } x \ P + p\text{-restrict } y \ Q) = p\text{-open-nbhd } p \ n \ y$

if $\neg P \ p$ **and** $Q \ p$

$\langle \text{proof} \rangle$

lemma *p-open-nbhd-antimono*:

$p\text{-open-nbhd } p \ n \ x \subseteq p\text{-open-nbhd } p \ m \ x$ **if** $m \leq n$

$\langle \text{proof} \rangle$

lemma *p-open-nbhds-open-subopen*:

$\text{generate-topology } (\text{local-}p\text{-open-nbhds } p) \ A =$

$(\forall a \in A. \exists n. p\text{-open-nbhd } p \ n \ a \subseteq A)$

$\langle \text{proof} \rangle$

lemma *p-restrict-p-open-set-mem-iff*:

$a \in A \longleftrightarrow p\text{-restrict } a \ P \in A$

if $\text{generate-topology } (\text{local-}p\text{-open-nbhds } p) \ A$ **and** $P \ p$

$\langle \text{proof} \rangle$

lemma *hausdorff*:

$\exists p \ n. (p\text{-open-nbhd } p \ n \ x) \cap (p\text{-open-nbhd } p \ n \ y) = \{\}$ **if** $x \neq y$

$\langle \text{proof} \rangle$

sublocale *t2-p-open-nbhds*: *t2-space generate-topology p-open-nbhds*

$\langle \text{proof} \rangle$

abbreviation *p-open-nbhds-tendsto* $\equiv t2\text{-}p\text{-open-nbhds.tendsto}$

abbreviation *p-open-nbhds-convergent* $\equiv t2\text{-}p\text{-open-nbhds.convergent}$

abbreviation *p-open-nbhds-LIMSEQ* $\equiv t2\text{-}p\text{-open-nbhds.LIMSEQ}$

lemma *locally-limD*: $\exists K. p\text{-limseq-condition } p \ X \ x \ n \ K$ **if** $p\text{-open-nbhds-LIMSEQ}$

$X \ x$

$\langle \text{proof} \rangle$

lemma *globally-limseq-imp-locally-limseq*:

$p\text{-limseq } p \ X \ x$ **if** $p\text{-open-nbhds-LIMSEQ } X \ x$

$\langle \text{proof} \rangle$

lemma *globally-limseq-iff-locally-limseq*: $p\text{-open-nbhds-LIMSEQ } X \ x = (\forall p. p\text{-limseq}$

$p \ X \ x)$

$\langle \text{proof} \rangle$

lemma *p-open-nbhds-LIMSEQ-eventually-p-constant*:

$p\text{-equal } p \ x \ c$

if $p\text{-open-nbhds-LIMSEQ } X \ x$ **and** $\forall_F k$ *in sequentially*. $p\text{-equal } p \ (X \ k) \ c$

$\langle \text{proof} \rangle$

lemma *p-open-nbhds-LIMSEQ-p-restrict*:
p-equal p x 0 **if** *p-open-nbhds-LIMSEQ* ($\lambda n. p\text{-restrict } (X\ n)\ P$) *x* **and** $\neg P\ p$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-p-open-nbhds-LIMSEQ-closed*:
 $x \in \text{global-depth-set } n$
if *p-open-nbhds-LIMSEQ* $X\ x$
and $\forall_F k$ *in sequentially*. $X\ k \in \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

end

context *global-p-equal-depth-div-inv-w-inj-index-consts*
begin

lemma *shift-p-depth-p-open-nbhd*:
 $\text{shift-p-depth } p\ n\ 'p\text{-open-nbhd } p\ m\ x = p\text{-open-nbhd } p\ (m + n)\ (\text{shift-p-depth } p\ n\ x)$
 $\langle \text{proof} \rangle$

lemma *shift-p-depth-p-open-set*:
 $\text{generate-topology } (\text{local-p-open-nbhds } p)\ (\text{shift-p-depth } p\ n\ 'A)$
if $\text{generate-topology } (\text{local-p-open-nbhds } p)\ A$
 $\langle \text{proof} \rangle$

end

locale *p-complete-global-p-equal-depth* = *global-p-equal-depth* +
assumes *complete*:
 $p\text{-cauchy } p\ X \implies p\text{-open-nbhds-convergent } (\lambda n. p\text{-restrict } (X\ n)\ ((=)\ p))$
begin

lemma *p-cauchyE*:
assumes *p-cauchy* $p\ X$
obtains x **where** *p-open-nbhds-LIMSEQ* ($\lambda n. p\text{-restrict } (X\ n)\ ((=)\ p)$) x
 $\langle \text{proof} \rangle$

end

Hensel's Lemma.

context *p-equality-div-inv*
begin

— A sequence to approach a root from within the ring of integers.

primrec *hensel-seq* :: $'a \Rightarrow 'b\ \text{poly} \Rightarrow 'b \Rightarrow \text{nat} \Rightarrow 'b$
where $\text{hensel-seq } p\ f\ x\ 0 = x \mid$
 $\text{hensel-seq } p\ f\ x\ (\text{Suc } n) = ($
 $\text{if } p\text{-equal } p\ (\text{poly } f\ (\text{hensel-seq } p\ f\ x\ n))\ 0 \text{ then } \text{hensel-seq } p\ f\ x\ n$
 else

$$\text{hensel-seq } p \text{ f } x \text{ n} - (\text{poly } f \text{ (hensel-seq } p \text{ f } x \text{ n)}) / (\text{poly } (\text{polyderiv } f) \text{ (hensel-seq } p \text{ f } x \text{ n)})$$

lemma *constant-hensel-seq-case:*

$k \geq n \implies \text{hensel-seq } p \text{ f } x \text{ k} = \text{hensel-seq } p \text{ f } x \text{ n}$
if $p\text{-equal } p \text{ (poly } f \text{ (hensel-seq } p \text{ f } x \text{ n)) } 0$
 <proof>

end

context *p-equality-depth-div-inv*
begin

— Context for inductive steps leading to Hensel's Lemma.

context

fixes $p :: 'a$ **and** $f :: 'b$ *poly* **and** $x :: 'b$ **and** $n :: \text{nat}$
assumes $f\text{-deg} : \text{degree } f > 0$
and $f\text{-depth} : \text{set } (\text{coeffs } f) \subseteq p\text{-depth-set } p \text{ } 0$
and $d\text{-hensel} : \text{hensel-seq } p \text{ f } x \text{ n} \in p\text{-depth-set } p \text{ } 0$
and $f : \neg p\text{-equal } p \text{ (poly } f \text{ (hensel-seq } p \text{ f } x \text{ n)) } 0$
and $\text{deriv-f} : \neg p\text{-equal } p \text{ (poly } (\text{polyderiv } f) \text{ (hensel-seq } p \text{ f } x \text{ n)) } 0$
and $df :$
 $p\text{-depth } p \text{ (poly } f \text{ (hensel-seq } p \text{ f } x \text{ n)) } >$
 $2 * p\text{-depth } p \text{ (poly } (\text{polyderiv } f) \text{ (hensel-seq } p \text{ f } x \text{ n))}$

begin

lemma *hensel-seq-step-1:*

$p\text{-depth } p \text{ (poly } f \text{ (hensel-seq } p \text{ f } x \text{ (Suc } n))) >$
 $p\text{-depth } p \text{ (poly } f \text{ ((hensel-seq } p \text{ f } x \text{ n)))}$
if $\text{next-f} : \neg p\text{-equal } p \text{ (poly } f \text{ (hensel-seq } p \text{ f } x \text{ (Suc } n))) } 0$
 <proof>

lemma *hensel-seq-step-2:*

$\neg p\text{-equal } p \text{ (poly } (\text{polyderiv } f) \text{ (hensel-seq } p \text{ f } x \text{ (Suc } n))) } 0$
 $p\text{-depth } p \text{ (poly } (\text{polyderiv } f) \text{ (hensel-seq } p \text{ f } x \text{ (Suc } n))) =$
 $p\text{-depth } p \text{ (poly } (\text{polyderiv } f) \text{ (hensel-seq } p \text{ f } x \text{ n))}$
 <proof>

end

— Context for Hensel's Lemma.

context

fixes $p :: 'a$ **and** $f :: 'b$ *poly* **and** $x :: 'b$
assumes $f\text{-deg} : \text{degree } f > 0$
and $f\text{-depth} : \text{set } (\text{coeffs } f) \subseteq p\text{-depth-set } p \text{ } 0$
and $d\text{-init} : x \in p\text{-depth-set } p \text{ } 0$
and $\text{init-deriv} :$
 $\neg p\text{-equal } p \text{ (poly } f \text{ } x) 0 \longrightarrow \neg p\text{-equal } p \text{ (poly } (\text{polyderiv } f) \text{ } x) 0$

and *d-init-deriv*:
 $\neg p\text{-equal } p \text{ (poly } f \text{ } x) \text{ } 0 \longrightarrow$
 $p\text{-depth } p \text{ (poly } f \text{ } x) > 2 * p\text{-depth } p \text{ (poly (polyderiv } f) \text{ } x)$
begin

lemma
hensel-seq-adelic-int : $\text{hensel-seq } p \text{ } f \text{ } x \text{ } n \in p\text{-depth-set } p \text{ } 0$ (**is** ?A)
and *hensel-seq-poly-depth* :
 $\neg p\text{-equal } p \text{ (poly } f \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n)) \text{ } 0 \longrightarrow$
 $p\text{-depth } p \text{ (poly } f \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n))} \geq p\text{-depth } p \text{ (poly } f \text{ } x) + \text{int } n$
(**is** ?B)
and *hensel-seq-deriv* :
 $\neg p\text{-equal } p \text{ (poly } f \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n)) \text{ } 0 \longrightarrow$
 $\neg p\text{-equal } p \text{ (poly (polyderiv } f) \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n)) \text{ } 0$
(**is** ?C)
and *hensel-seq-deriv-depth*:
 $\neg p\text{-equal } p \text{ (poly } f \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n)) \text{ } 0 \longrightarrow$
 $p\text{-depth } p \text{ (poly (polyderiv } f) \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n))} = p\text{-depth } p \text{ (poly (polyderiv$
 $f) \text{ } x)$
(**is** ?D)
and *hensel-seq-depth-ineq*:
 $\neg p\text{-equal } p \text{ (poly } f \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n)) \text{ } 0 \longrightarrow$
 $p\text{-depth } p \text{ (poly } f \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n))} >$
 $2 * p\text{-depth } p \text{ (poly (polyderiv } f) \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n))}$
(**is** ?E)
 $\langle \text{proof} \rangle$

lemma *hensel-seq-cauchy*: $p\text{-cauchy } p \text{ (hensel-seq } p \text{ } f \text{ } x)$
 $\langle \text{proof} \rangle$

lemma *p-limseq-poly-hensel-seq*: $p\text{-limseq } p \text{ (}\lambda n. \text{poly } f \text{ (hensel-seq } p \text{ } f \text{ } x \text{ } n)) \text{ } 0$
 $\langle \text{proof} \rangle$

end

end

locale *p-complete-global-p-equal-depth-div-inv* =
 $\text{global-p-equal-depth-div-inv} + p\text{-complete-global-p-equal-depth}$
begin

lemma *hensel*:
fixes $p :: 'a$ **and** $f :: 'b \text{ poly}$ **and** $x :: 'b$
assumes $\text{degree } f > 0$
and $\text{set (coeffs } f) \subseteq p\text{-depth-set } p \text{ } 0$
and $x \in p\text{-depth-set } p \text{ } 0$
and
 $\neg p\text{-equal } p \text{ (poly } f \text{ } x) \text{ } 0 \longrightarrow \neg p\text{-equal } p \text{ (poly (polyderiv } f) \text{ } x) \text{ } 0$

and
 $\neg p\text{-equal } p \text{ (poly } f \text{) } 0 \longrightarrow$
 $p\text{-depth } p \text{ (poly } f \text{) } > 2 * p\text{-depth } p \text{ (poly (polyderiv } f \text{) } x)$
obtains z **where** $z \in p\text{-depth-set } p \text{ } 0$ **and** $p\text{-equal } p \text{ (poly } f \text{) } z \text{ } 0$
 $\langle \text{proof} \rangle$
end

3.7.2 Globally

We use bounded depth to create a global metric.

context $p\text{-equality-depth}$
begin

definition $bdd\text{-global-depth} :: 'b \Rightarrow nat$
where
 $bdd\text{-global-depth } x =$
 $(\text{if globally-}p\text{-equal } x \text{ } 0 \text{ then } 0 \text{ else } \text{Inf } \{nat \text{ (} p\text{-depth } p \text{ } x + 1) \mid p. \neg p\text{-equal } p$
 $x \text{ } 0\})$
 — The plus-one is to distinguish depth-0 from negative depth.

lemma $bdd\text{-global-depth-0[simp]: globally-}p\text{-equal } x \text{ } 0 \implies bdd\text{-global-depth } x = 0$
 $\langle \text{proof} \rangle$

lemma $bdd\text{-global-depthD:}$
 $bdd\text{-global-depth } x = \text{Inf } \{nat \text{ (} p\text{-depth } p \text{ } x + 1) \mid p. \neg p\text{-equal } p \text{ } x \text{ } 0\}$
if $\neg \text{globally-}p\text{-equal } x \text{ } 0$
 $\langle \text{proof} \rangle$

lemma $bdd\text{-global-depthD-as-image:}$
 $bdd\text{-global-depth } x = (\text{INF } p \in \{p. \neg p\text{-equal } p \text{ } x \text{ } 0\}. nat \text{ (} p\text{-depth } p \text{ } x + 1))$
if $\neg \text{globally-}p\text{-equal } x \text{ } 0$
 $\langle \text{proof} \rangle$

lemma $bdd\text{-global-depth-cong:}$
 $bdd\text{-global-depth } x = bdd\text{-global-depth } y$ **if** $\text{globally-}p\text{-equal } x \text{ } y$
 $\langle \text{proof} \rangle$

lemma $bdd\text{-global-depth-le:}$
 $bdd\text{-global-depth } x \leq nat \text{ (} p\text{-depth } p \text{ } x + 1)$ **if** $\neg p\text{-equal } p \text{ } x \text{ } 0$
 $\langle \text{proof} \rangle$

lemma $bdd\text{-global-depth-greatest:}$
 $bdd\text{-global-depth } x \geq B$
if $\neg p\text{-equal } p \text{ } x \text{ } 0$
and $\forall q. \neg p\text{-equal } q \text{ } x \text{ } 0 \longrightarrow nat \text{ (} p\text{-depth } q \text{ } x + 1) \geq B$
 $\langle \text{proof} \rangle$

lemma $bdd\text{-global-depth-witness:}$

assumes $\neg \text{globally-p-equal } x \ 0$
obtains p **where** $\neg p\text{-equal } p \ x \ 0$ **and** $\text{bdd-global-depth } x = \text{nat } (p\text{-depth } p \ x + 1)$
 $\langle \text{proof} \rangle$

lemma *bdd-global-depth-eq-0*:
 $\text{bdd-global-depth } x = 0$ **if** $p\text{-depth } p \ x < 0$
 $\langle \text{proof} \rangle$

lemma *bdd-global-depth-eq-0-iff*:
 $\text{bdd-global-depth } x = 0 \iff (\exists p. p\text{-depth } p \ x < 0) \vee (\text{globally-p-equal } x \ 0)$
(is $?L = ?R)$
 $\langle \text{proof} \rangle$

lemma *bdd-global-depth-diff*: $\text{bdd-global-depth } (x - y) = \text{bdd-global-depth } (y - x)$
 $\langle \text{proof} \rangle$

lemma *bdd-global-depth-uminus*: $\text{bdd-global-depth } (-x) = \text{bdd-global-depth } x$
 $\langle \text{proof} \rangle$

lemma *bdd-global-depth-nonarch*:
 $\text{bdd-global-depth } (x + y) \geq \min (\text{bdd-global-depth } x) (\text{bdd-global-depth } y)$
if $\neg \text{globally-p-equal } (x + y) \ 0$
 $\langle \text{proof} \rangle$

definition *bdd-global-dist* :: $'b \Rightarrow 'b \Rightarrow \text{real}$ **where**
 $\text{bdd-global-dist } x \ y =$
 $(\text{if } \text{globally-p-equal } x \ y \text{ then } 0 \text{ else } \text{inverse } (2 \wedge \text{bdd-global-depth } (x - y)))$

lemma *bdd-global-dist-eqD [simp]*: $\text{bdd-global-dist } x \ y = 0$ **if** $\text{globally-p-equal } x \ y$
 $\langle \text{proof} \rangle$

lemma *bdd-global-distD*:
 $\text{bdd-global-dist } x \ y = \text{inverse } (2 \wedge \text{bdd-global-depth } (x - y))$ **if** $\neg \text{globally-p-equal } x \ y$
 $\langle \text{proof} \rangle$

lemma *bdd-global-dist-cong*:
 $\text{bdd-global-dist } w \ x = \text{bdd-global-dist } y \ z$
if $\text{globally-p-equal } w \ y$ **and** $\text{globally-p-equal } x \ z$
 $\langle \text{proof} \rangle$

lemma *bdd-global-dist-nonneg*: $\text{bdd-global-dist } x \ y \geq 0$
 $\langle \text{proof} \rangle$

lemma *bdd-global-dist-bdd*: $\text{bdd-global-dist } x \ y \leq 1$
 $\langle \text{proof} \rangle$

lemma *bdd-global-dist-lessI*:

bdd-global-dist $x\ y < e$

if *pos*: $e > 0$

and *depth*:

$\bigwedge p. \neg p\text{-equal}\ p\ x\ y \implies$

$p\text{-depth}\ p\ (x - y) \geq \lfloor \log 2\ (\text{inverse } e) \rfloor$

<proof>

lemma *bdd-global-dist-less-imp*:

$p\text{-depth}\ p\ (x - y) \geq \lfloor \log 2\ (\text{inverse } e) \rfloor$

if $e > 0$ **and** $e \leq 1$ **and** $\neg p\text{-equal}\ p\ x\ y$ **and** *bdd-global-dist* $x\ y < e$

<proof>

lemma *bdd-global-dist-less-pow2-iff*:

bdd-global-dist $x\ y < \text{inverse}\ (2 \wedge n) \iff$

$(\forall p. \neg p\text{-equal}\ p\ x\ y \implies p\text{-depth}\ p\ (x - y) \geq \text{int } n)$

<proof>

lemma *bdd-global-dist-sym*: *bdd-global-dist* $x\ y = \text{bdd-global-dist}\ y\ x$

<proof>

lemma *bdd-global-dist-conv0*: *bdd-global-dist* $x\ y = \text{bdd-global-dist}\ (x - y)\ 0$

<proof>

lemma *bdd-global-dist-conv0'*: *bdd-global-dist* $x\ y = \text{bdd-global-dist}\ 0\ (x - y)$

<proof>

lemma *bdd-global-dist-nonarch*:

bdd-global-dist $x\ y \leq \max\ (\text{bdd-global-dist}\ x\ z)\ (\text{bdd-global-dist}\ y\ z)$

<proof>

lemma *bdd-global-dist-ball-multicentre*:

$\{z. \text{bdd-global-dist}\ y\ z < e\} = \{z. \text{bdd-global-dist}\ x\ z < e\}$

if *bdd-global-dist* $x\ y < e$

<proof>

lemma *bdd-global-dist-ball-at-0*:

$\{z. \text{bdd-global-dist}\ 0\ z < \text{inverse}\ (2 \wedge n)\} = \text{global-depth-set}\ (\text{int } n)$

<proof>

lemma *bdd-global-dist-ball-UNIV*:

$\{z. \text{bdd-global-dist}\ x\ z < e\} = \text{UNIV}$ **if** $e > 1$

<proof>

lemma *bdd-global-dist-ball-at-0-normalize*:

$\{z. \text{bdd-global-dist}\ 0\ z < e\} = \text{global-depth-set}\ (\lfloor \log 2\ (\text{inverse } e) \rfloor)$

if $e > 0$ **and** $e \leq 1$

<proof>

lemma *bdd-global-dist-ball-translate:*

$$\{z. \text{bdd-global-dist } x \ z < e\} = x + o \ \{z. \text{bdd-global-dist } 0 \ z < e\}$$

<proof>

lemma *bdd-global-dist-left-translate-continuous:*

$$\text{bdd-global-dist } (x + y) \ (x + z) < e \text{ if } \text{bdd-global-dist } y \ z < e$$

<proof>

lemma *bdd-global-dist-right-translate-continuous:*

$$\text{bdd-global-dist } (y + x) \ (z + x) < e \text{ if } \text{bdd-global-dist } y \ z < e$$

<proof>

definition *global-cauchy-condition ::*

$$(nat \Rightarrow 'b) \Rightarrow nat \Rightarrow nat \Rightarrow bool$$

where

$$\text{global-cauchy-condition } X \ n \ K =$$

$$(\forall k1 \geq K. \forall k2 \geq K. \forall p.$$

$$\neg p\text{-equal } p \ (X \ k1) \ (X \ k2) \longrightarrow nat \ (p\text{-depth } p \ (X \ k1 - X \ k2)) > n$$

)

lemma *global-cauchy-conditionI:*

$$\text{global-cauchy-condition } X \ n \ K$$

if

$$\bigwedge p \ k \ k'. \ k \geq K \implies k' \geq K \implies$$

$$\neg p\text{-equal } p \ (X \ k) \ (X \ k') \implies nat \ (p\text{-depth } p \ (X \ k - X \ k')) > n$$

<proof>

lemma *global-cauchy-conditionD:*

$$nat \ (p\text{-depth } p \ (X \ k - X \ k')) > n$$

$$\text{if } \text{global-cauchy-condition } X \ n \ K \text{ and } k \geq K \text{ and } k' \geq K$$

$$\text{and } \neg p\text{-equal } p \ (X \ k) \ (X \ k')$$

<proof>

lemma *global-cauchy-condition-mono-seq-tail:*

$$\text{global-cauchy-condition } X \ n \ K \implies \text{global-cauchy-condition } X \ n \ K'$$

$$\text{if } K' \geq K$$

<proof>

lemma *global-cauchy-condition-mono-depth:*

$$\text{global-cauchy-condition } X \ n \ K \implies \text{global-cauchy-condition } X \ m \ K$$

$$\text{if } m \leq n$$

<proof>

abbreviation *globally-cauchy* $X \equiv (\forall n. \exists K. \text{global-cauchy-condition } X \ n \ K)$

lemma *global-cauchy-condition-LEAST:*

$$\text{global-cauchy-condition } X \ n \ (\text{LEAST } K. \text{global-cauchy-condition } X \ n \ K) \text{ if globally-cauchy } X$$

<proof>

lemma *global-cauchy-condition-LEAST-mono*:
 $(LEAST\ K.\ global\text{-}cauchy\text{-}condition\ X\ m\ K) \leq (LEAST\ K.\ global\text{-}cauchy\text{-}condition\ X\ n\ K)$
if *globally-cauchy* X **and** $m \leq n$
 $\langle proof \rangle$

definition *bdd-global-uniformity* :: $('b \times 'b)$ filter
where *bdd-global-uniformity* = $(INF\ e \in \{0 < ..\}. principal\ \{(x, y). bdd\text{-}global\text{-}dist\ x\ y < e\})$

end

context *global-p-equal-depth*
begin

sublocale *metric-space-by-bdd-depth*:
metric-space *bdd-global-dist* *bdd-global-uniformity*
 $\lambda U. \forall x \in U.$
eventually $(\lambda(x', y). x' = x \longrightarrow y \in U)$ *bdd-global-uniformity*
 $\langle proof \rangle$

lemma *nonneg-global-depth-set-LIMSEQ-closed*:
 $x \in global\text{-}depth\text{-}set\ (int\ n)$
if *lim*: *metric-space-by-bdd-depth.LIMSEQ* $X\ x$
and *seq*: $\forall_F k\ in\ sequentially. X\ k \in global\text{-}depth\text{-}set\ (int\ n)$
 $\langle proof \rangle$

lemma *globally-cauchy-vs-bdd-metric-Cauchy*:
globally-cauchy $X = metric\text{-}space\text{-}by\text{-}bdd\text{-}depth.Cauchy\ X$ **for** $X :: nat \Rightarrow 'b$
 $\langle proof \rangle$

end

context *global-p-equal-depth-no-zero-divisors*
begin

lemma *bdd-global-dist-bdd-mult-continuous*:
 $bdd\text{-}global\text{-}dist\ (w * x)\ (w * y) < e$
if *pos* : $e > 0$
and *bdd* : $\forall p. \neg p\text{-equal}\ p\ w\ 0 \longrightarrow p\text{-depth}\ p\ w \geq n$
and *input*: $bdd\text{-}global\text{-}dist\ x\ y < min\ (e * 2\ powi\ (n - 1))\ 1$
 $\langle proof \rangle$

lemma *bdd-global-dist-limseq-bdd-mult*:
metric-space-by-bdd-depth.LIMSEQ $(\lambda k. w * X\ k)\ (w * x)$
if *bdd*: $\forall p. \neg p\text{-equal}\ p\ w\ 0 \longrightarrow p\text{-depth}\ p\ w \geq n$
and *seq*: *metric-space-by-bdd-depth.LIMSEQ* $X\ x$
 $\langle proof \rangle$

end

context *global-p-equal-depth-div-inv-w-inj-index-consts*
begin

lemma *bdd-global-dist-limseq-global-uniformizer-powi-mult:*
metric-space-by-bdd-depth.LIMSEQ
 $(\lambda k. \text{global-uniformizer powi } n * X \ k) \ (\text{global-uniformizer powi } n * x)$
if *metric-space-by-bdd-depth.LIMSEQ* $X \ x$
 $\langle \text{proof} \rangle$

lemma *global-depth-set-LIMSEQ-closed:*
 $x \in \text{global-depth-set } n$
if *lim* : *metric-space-by-bdd-depth.LIMSEQ* $X \ x$
and *depth*: $\forall_F \ k \text{ in sequentially. } X \ k \in \text{global-depth-set } n$
 $\langle \text{proof} \rangle$

end

3.8 Notation

consts
 $p\text{-equal} :: 'a \Rightarrow 'b \Rightarrow 'b \Rightarrow \text{bool}$
 $globally\text{-}p\text{-equal} ::$
 $'b \Rightarrow 'b \Rightarrow \text{bool}$
 $((-/ \simeq_{\forall p} / -) [51, 51] 50)$
 $p\text{-restrict} :: 'b \Rightarrow ('a \Rightarrow \text{bool}) \Rightarrow 'b \ (\text{infixl } prestrict \ 70)$
 $p\text{-depth} :: 'a \Rightarrow 'b \Rightarrow \text{int}$
 $p\text{-depth-set} :: 'a \Rightarrow \text{int} \Rightarrow 'b \text{ set } (\mathcal{P}_{@-})$
 $global\text{-}depth\text{-set} :: \text{int} \Rightarrow 'b \text{ set } (\mathcal{P}_{\forall p-})$
 $global\text{-}unfrmzr\text{-}pows :: ('a \Rightarrow 'c) \Rightarrow 'b$

abbreviation $p\text{-equal-notation} :: 'b \Rightarrow 'a \Rightarrow 'b \Rightarrow \text{bool}$
 $((-/ \simeq / -) [51, 51, 51] 50)$
where $p\text{-equal-notation } x \ p \ y \equiv p\text{-equal } p \ x \ y$

abbreviation $p\text{-nequal-notation} :: 'b \Rightarrow 'a \Rightarrow 'b \Rightarrow \text{bool}$
 $((-/ \neg \simeq / -) [51, 51, 51] 50)$
where $p\text{-nequal-notation } x \ p \ y \equiv \neg \ p\text{-equal } p \ x \ y$

abbreviation $p\text{-depth-notation} :: 'b \Rightarrow 'a \Rightarrow \text{int}$
 $((-\circ-) [51, 51] 50)$
where $p\text{-depth-notation } x \ p \equiv p\text{-depth } p \ x$

abbreviation $p\text{-depth-set-0-notation} :: 'a \Rightarrow 'b \text{ set } (\mathcal{O}_{@-})$
where $p\text{-depth-set-0-notation } p \equiv p\text{-depth-set } p \ 0$

abbreviation $p\text{-depth-set-1-notation} :: 'a \Rightarrow 'b \text{ set } (\mathcal{P}_{@-})$

where $p\text{-depth-set-1-notation } p \equiv p\text{-depth-set } p \ 1$
abbreviation $global\text{-depth-set-0-notation} :: 'b \text{ set } (\mathcal{O}_{\forall p})$
where $global\text{-depth-set-0-notation} \equiv global\text{-depth-set } 0$
abbreviation $global\text{-depth-set-1-notation} :: 'b \text{ set } (\mathcal{P}_{\forall p})$
where $global\text{-depth-set-1-notation} \equiv global\text{-depth-set } 1$
abbreviation $\mathfrak{p} \equiv global\text{-unfrmzr-pows}$
end

theory *FLS-Prime-Equiv-Depth*

imports
Parametric-Equiv-Depth
HOL.Equiv-Relations
HOL-Computational-Algebra.Formal-Laurent-Series

begin

unbundle *fps-syntax*

4 Equivalence of sequences of formal Laurent series relative to divisibility of partial sums by a fixed prime

4.1 Preliminaries

4.1.1 Lift/transfer of equivalence relations.

lemma *reflp-transfer*:
 $reflp \ r \implies reflp \ (\lambda x \ y. \ r \ (f \ x) \ (f \ y))$
<proof>

lemma *symp-transfer*:
 $symp \ r \implies symp \ (\lambda x \ y. \ r \ (f \ x) \ (f \ y))$
<proof>

lemma *transp-transfer*:
 $transp \ r \implies transp \ (\lambda x \ y. \ r \ (f \ x) \ (f \ y))$
<proof>

lemma *equivp-transfer*:
 $equivp \ r \implies equivp \ (\lambda x \ y. \ r \ (f \ x) \ (f \ y))$
<proof>

4.1.2 An additional fact about products/sums

lemma (in *comm-monoid-set*) *atLeastAtMost-int-shift-bounds*:
 $F\ g\ \{m..n\} = F\ (g \circ \text{plus } m \circ \text{int})\ \{..nat\ (n - m)\}$
 if $m \leq n$
 for $m\ n :: \text{int}$
 <proof>

4.1.3 An additional fact about algebraic powers of functions

lemma *pow-fun-apply*: $(f \wedge n)\ x = (f\ x) \wedge n$
 <proof>

4.1.4 Some additional facts about formal power and Laurent series

lemma *fps-shift-fps-mult-by-fps-const*:
 fixes $c :: 'a::\{\text{comm-monoid-add,mult-zero}\}$
 shows $\text{fps-shift } n\ (\text{fps-const } c * f) = \text{fps-const } c * \text{fps-shift } n\ f$
 and $\text{fps-shift } n\ (f * \text{fps-const } c) = \text{fps-shift } n\ f * \text{fps-const } c$
 <proof>

lemma *fps-shift-mult-Suc*:
 fixes $f\ g :: 'a::\{\text{comm-monoid-add,mult-zero}\}\ \text{fps}$
 shows $\text{fps-shift } (\text{Suc } n)\ (f * g)$
 $= \text{fps-const } (f\$0) * \text{fps-shift } (\text{Suc } n)\ g + \text{fps-shift } n\ (\text{fps-shift } 1\ f * g)$
 and $\text{fps-shift } (\text{Suc } n)\ (f * g)$
 $= \text{fps-shift } n\ (f * \text{fps-shift } 1\ g) + \text{fps-shift } (\text{Suc } n)\ f * \text{fps-const } (g\$0)$
 <proof>

lemma *fls-subdegree-minus'*:
 $\text{fls-subdegree } (f - g) = \text{fls-subdegree } f$
 if $f \neq 0$ and $\text{fls-subdegree } g > \text{fls-subdegree } f$
 <proof>

lemma *fls-regpart-times-fps-X*:
 fixes $f :: 'a::\{\text{comm-monoid-add,mult-zero,monoid-mult}\}\ \text{fls}$
 assumes $\text{fls-subdegree } f \geq 0$
 shows $\text{fls-regpart } f * \text{fps-X} = \text{fls-regpart } (\text{fls-shift } (-1)\ f)$
 and $\text{fps-X} * \text{fls-regpart } f = \text{fls-regpart } (\text{fls-shift } (-1)\ f)$
 <proof>

lemma *fls-regpart-times-fps-X-power*:
 fixes $f :: 'a::\text{semiring-1}\ \text{fls}$
 assumes $\text{fls-subdegree } f \geq 0$
 shows $\text{fls-regpart } f * \text{fps-X} \wedge n = \text{fls-regpart } (\text{fls-shift } (-n)\ f)$
 and $\text{fps-X} \wedge n * \text{fls-regpart } f = \text{fls-regpart } (\text{fls-shift } (-n)\ f)$
 <proof>

lemma *fls-base-factor-uminus*: $\text{fls-base-factor } (-f) = -\ \text{fls-base-factor } f$

$\langle \text{proof} \rangle$

4.2 Partial evaluation of formal power series

Make a degree- n polynomial out of a formal power series.

definition *poly-of-fps*

where $\text{poly-of-fps } f \ n = \text{Abs-poly } (\lambda i. \text{ if } i \leq n \text{ then } f\$i \text{ else } 0)$

lemma *if-then-MOST-zero*:

$\forall_\infty x. (\text{if } P \ x \text{ then } f \ x \text{ else } 0) = 0$

if $\forall_\infty x. \neg P \ x$

$\langle \text{proof} \rangle$

lemma *truncated-fps-eventually-zero*:

$(\lambda i. \text{ if } i \leq n \text{ then } f\$i \text{ else } 0) \in \{g. \forall_\infty n. g \ n = 0\}$

$\langle \text{proof} \rangle$

lemma *coeff-poly-of-fps[simp]*:

$\text{coeff } (\text{poly-of-fps } f \ n) \ i = (\text{if } i \leq n \text{ then } f\$i \text{ else } 0)$

$\langle \text{proof} \rangle$

lemma *poly-of-fps-0[simp]*: $\text{poly-of-fps } 0 \ n = 0$

$\langle \text{proof} \rangle$

lemma *poly-of-fps-0th-conv-pCons[simp]*: $\text{poly-of-fps } f \ 0 = [:f\$0:]$

$\langle \text{proof} \rangle$

lemma *poly-of-fps-degree*: $\text{degree } (\text{poly-of-fps } f \ n) \leq n$

$\langle \text{proof} \rangle$

lemma *poly-of-fps-Suc*:

$\text{poly-of-fps } f \ (\text{Suc } n) = \text{poly-of-fps } f \ n + \text{monom } (f \ \$ \ \text{Suc } n) \ (\text{Suc } n)$

$\langle \text{proof} \rangle$

lemma *poly-of-fps-Suc-nth-eq0*:

$\text{poly-of-fps } f \ (\text{Suc } n) = \text{poly-of-fps } f \ n \text{ if } f \ \$ \ \text{Suc } n = 0$

$\langle \text{proof} \rangle$

lemma *poly-of-fps-Suc-conv-pCons-shift*:

$\text{poly-of-fps } f \ (\text{Suc } n) = \text{pCons } (f\$0) \ (\text{poly-of-fps } (\text{fps-shift } 1 \ f) \ n)$

$\langle \text{proof} \rangle$

abbreviation *fps-parteval* $f \ x \ n \equiv \text{poly } (\text{poly-of-fps } f \ n) \ x$

abbreviation *fls-base-parteval* $f \equiv \text{fps-parteval } (\text{fls-base-factor-to-fps } f)$

lemma *fps-parteval-0[simp]*: $\text{fps-parteval } 0 \ x \ n = 0$

$\langle \text{proof} \rangle$

lemma *fps-parteval-0th[simp]*: $\text{fps-parteval } f \ x \ 0 = f \ \$ \ 0$

$\langle \text{proof} \rangle$

lemma *fps-parteval-Suc*:

$\text{fps-parteval } f \ x \ (\text{Suc } n) = \text{fps-parteval } f \ x \ n + (f \ \$ \ \text{Suc } n) * x^{\wedge} \text{Suc } n$

for $f :: 'a::\text{comm-semiring-1}$ *fps*

$\langle \text{proof} \rangle$

lemma *fps-parteval-Suc-cons*:

$\text{fps-parteval } f \ x \ (\text{Suc } n) = f \$ 0 + x * \text{fps-parteval } (\text{fps-shift } 1 \ f) \ x \ n$

$\langle \text{proof} \rangle$

lemma *fps-parteval-at-0[simp]*: $\text{fps-parteval } f \ 0 \ n = f \ \$ \ 0$

$\langle \text{proof} \rangle$

lemma *fps-parteval-conv-sum*:

$\text{fps-parteval } f \ x \ n = (\sum k \leq n. f \$ k * x^{\wedge} k)$ **for** $x :: 'a::\text{comm-semiring-1}$

$\langle \text{proof} \rangle$

lemma *fls-base-parteval-conv-sum*:

$\text{fls-base-parteval } f \ x \ n = (\sum k \leq n. f \$ \$ (\text{fls-subdegree } f + \text{int } k) * x^{\wedge} k)$

for $f :: 'a::\text{comm-semiring-1}$ *fls* **and** $x :: 'a$ **and** $n :: \text{nat}$

$\langle \text{proof} \rangle$

lemma *fps-parteval-fps-const[simp]*:

$\text{fps-parteval } (\text{fps-const } c) \ x \ n = c$

$\langle \text{proof} \rangle$

lemma *fps-parteval-1[simp]*: $\text{fps-parteval } 1 \ x \ n = 1$

$\langle \text{proof} \rangle$

lemma *fps-parteval-mult-by-fps-const*:

$\text{fps-parteval } (\text{fps-const } c * f) \ x \ n = c * \text{fps-parteval } f \ x \ n$

$\langle \text{proof} \rangle$

lemma *fps-parteval-plus[simp]*:

$\text{fps-parteval } (f + g) \ x \ n = \text{fps-parteval } f \ x \ n + \text{fps-parteval } g \ x \ n$

$\langle \text{proof} \rangle$

lemma *fps-parteval-uminus[simp]*:

$\text{fps-parteval } (-f) \ x \ n = - \text{fps-parteval } f \ x \ n$ **for** $f :: 'a::\text{comm-ring}$ *fps*

$\langle \text{proof} \rangle$

lemma *fps-parteval-minus[simp]*:

$\text{fps-parteval } (f - g) \ x \ n = \text{fps-parteval } f \ x \ n - \text{fps-parteval } g \ x \ n$

for $f :: 'a::\text{comm-ring}$ *fps*

$\langle \text{proof} \rangle$

lemma *fps-parteval-mult-fps-X*:

fixes $f :: 'a::\text{comm-semiring-1}$ *fps*

shows $\text{fps-parteval } (f * \text{fps-X}) \ x \ (\text{Suc } n) = x * \text{fps-parteval } f \ x \ n$
and $\text{fps-parteval } (\text{fps-X} * f) \ x \ (\text{Suc } n) = x * \text{fps-parteval } f \ x \ n$
 $\langle \text{proof} \rangle$

lemma $\text{fps-parteval-mult-fps-X-power}$:
fixes $f :: 'a::\text{comm-semiring-1} \ \text{fps}$
assumes $n \geq m$
shows $\text{fps-parteval } (f * \text{fps-X}^{\wedge m}) \ x \ n = x^{\wedge m} * \text{fps-parteval } f \ x \ (n - m)$
and $\text{fps-parteval } (\text{fps-X}^{\wedge m} * f) \ x \ n = x^{\wedge m} * \text{fps-parteval } f \ x \ (n - m)$
 $\langle \text{proof} \rangle$

lemma $\text{fps-parteval-mult-right}$:
 $\text{fps-parteval } (f * g) \ x \ n =$
 $\text{fps-parteval } (\text{Abs-fps } (\lambda i. f \$ i * \text{fps-parteval } g \ x \ (n - i))) \ x \ n$
 $\langle \text{proof} \rangle$

lemma $\text{fps-parteval-mult-left}$:
 $\text{fps-parteval } (f * g) \ x \ n =$
 $\text{fps-parteval } (\text{Abs-fps } (\lambda i. \text{fps-parteval } f \ x \ (n - i) * g \$ i)) \ x \ n$
 $\langle \text{proof} \rangle$

lemma $\text{fps-parteval-mult-conv-sum}$:
fixes $f \ g :: 'a::\text{comm-semiring-1} \ \text{fps}$
shows $\text{fps-parteval } (f * g) \ x \ n = (\sum k \leq n. f \$ k * \text{fps-parteval } g \ x \ (n - k) * x^{\wedge k})$
and $\text{fps-parteval } (f * g) \ x \ n = (\sum k \leq n. \text{fps-parteval } f \ x \ (n - k) * g \$ k * x^{\wedge k})$
 $\langle \text{proof} \rangle$

lemma $\text{fps-parteval-fls-base-factor-to-fps-diff}$:
 $\text{fls-base-parteval } (f - g) \ x \ n = \text{fls-base-parteval } f \ x \ n$
if $f \neq 0$ **and** $\text{fls-subdegree } g > \text{fls-subdegree } f + (\text{int } n)$
for $f \ g :: 'a::\text{comm-ring-1} \ \text{fls}$
 $\langle \text{proof} \rangle$

4.3 Vanishing of formal power series relative to divisibility of all partial sums

definition $\text{fps-pvanishes} :: 'a::\text{comm-semiring-1} \Rightarrow 'a \ \text{fps} \Rightarrow \text{bool}$
where $\text{fps-pvanishes } p \ f \equiv (\forall n::\text{nat}. p^{\wedge} \text{Suc } n \ \text{dvd} \ (\text{fps-parteval } f \ p \ n))$

lemma fps-pvanishesD : $\text{fps-pvanishes } p \ f \Longrightarrow p^{\wedge} \text{Suc } n \ \text{dvd} \ (\text{fps-parteval } f \ p \ n)$
 $\langle \text{proof} \rangle$

lemma fps-pvanishesI :
 $\text{fps-pvanishes } p \ f$ **if** $\bigwedge n::\text{nat}. p^{\wedge} \text{Suc } n \ \text{dvd} \ (\text{fps-parteval } f \ p \ n)$
 $\langle \text{proof} \rangle$

lemma $\text{fps-pvanishes-0[simp]}$: $\text{fps-pvanishes } p \ 0$
 $\langle \text{proof} \rangle$

lemma *fps-pvanishes-at-0[simp]*: *fps-pvanishes 0 f = (f\$0 = 0)*
 ⟨proof⟩

lemma *fps-pvanishes-at-unit*: *fps-pvanishes p f if is-unit p*
 ⟨proof⟩

lemma *fps-pvanishes-one-iff*:
fps-pvanishes p 1 = is-unit p
 ⟨proof⟩

lemma *fps-pvanishes-uminus[simp]*:
fps-pvanishes p (-f) if fps-pvanishes p f for f :: 'a::comm-ring-1 fps
 ⟨proof⟩

lemma *fps-pvanishes-add[simp]*:
fps-pvanishes p (f + g) if fps-pvanishes p f and fps-pvanishes p g
for f :: 'a::comm-semiring-1 fps
 ⟨proof⟩

lemma *fps-pvanishes-minus[simp]*:
fps-pvanishes p (f - g) if fps-pvanishes p f and fps-pvanishes p g
for f :: 'a::comm-ring-1 fps
 ⟨proof⟩

lemma *fps-pvanishes-mult*: *fps-pvanishes p (f * g) if fps-pvanishes p g*
 ⟨proof⟩

lemma *fps-pvanishes-mult-fps-X-iff*:
assumes (p::'a::idom) ≠ 0
*shows fps-pvanishes p f ⟷ fps-pvanishes p (f * fps-X)*
*and fps-pvanishes p f ⟷ fps-pvanishes p (fps-X * f)*
 ⟨proof⟩

lemma *fps-pvanishes-mult-fps-X-power-iff*:
assumes (p::'a::idom) ≠ 0
*shows fps-pvanishes p f ⟷ fps-pvanishes p (f * fps-X ^ n)*
*and fps-pvanishes p f ⟷ fps-pvanishes p (fps-X ^ n * f)*
 ⟨proof⟩

4.4 Equivalence and depth of formal Laurent series relative to a prime

4.4.1 Definition of equivalence and basic facts

abbreviation *fps-primevanishes p* $\equiv$ *fps-pvanishes (Rep-prime p)*

overloading

p-equal-fls $\equiv$
p-equal :: 'a::nontrivial-factorial-comm-ring prime $\Rightarrow$
'a fls $\Rightarrow$ *'a fls* $\Rightarrow$ *bool*

begin

definition *p-equal-fls* ::

'a::nontrivial-factorial-comm-ring prime $\Rightarrow$

'a fls $\Rightarrow$ *'a fls* $\Rightarrow$ *bool*

where

p-equal-fls p f g $\equiv$ *fps-primevanishes p (fls-base-factor-to-fps (f - g))*

end

context

fixes *p* :: *'a* :: *nontrivial-factorial-comm-ring prime*

begin

lemma *p-equal-flsD*:

f $\simeq_p$ *g* $\Longrightarrow$ *fps-primevanishes p (fls-base-factor-to-fps (f - g))*

for *f g* :: *'a fls*

<proof>

lemma *p-equal-fls-imp-p-dvd*: *p dvd ((f - g) \$\$ (fls-subdegree (f - g)))*

if *f* $\simeq_p$ *g*

for *f g* :: *'a fls*

<proof>

lemma *p-equal-flsI*:

fps-primevanishes p (fls-base-factor-to-fps (f - g)) $\Longrightarrow$ *f* $\simeq_p$ *g*

for *f g* :: *'a fls*

<proof>

lemma *p-equal-fls-conv-p-equal-0*:

f $\simeq_p$ *g* $\longleftrightarrow$ *(f - g)* $\simeq_p$ *0*

for *f g* :: *'a fls*

<proof>

lemma *p-equal-fls-refl*: *f* $\simeq_p$ *f* **for** *f* :: *'a fls*

<proof>

lemma *p-equal-fls-sym*:

f $\simeq_p$ *g* $\Longrightarrow$ *g* $\simeq_p$ *f* **for** *f g* :: *'a fls*

<proof>

lemma *p-equal-fls-shift-iff*:

f $\simeq_p$ *g* $\longleftrightarrow$ *(fls-shift n f)* $\simeq_p$ *(fls-shift n g)*

for *f g* :: *'a fls*

<proof>

lemma *p-equal-fls-extended-sum-iff*:

f $\simeq_p$ *g* $\longleftrightarrow$

$(\forall n. (\text{Rep-prime } p) \wedge \text{Suc } n \text{ dvd}$

$(\sum_{k \leq n}. (f - g) \$ (N + \text{int } k) * (\text{Rep-prime } p) \wedge k))$
if $N \leq \min (\text{fls-subdegree } f) (\text{fls-subdegree } g)$
 <proof>

lemma *p-equal-fls-extended-intsum-iff*:

$f \simeq_p g \longleftrightarrow$
 $(\forall n \geq N. (\text{Rep-prime } p) \wedge \text{Suc } (\text{nat } (n - N)) \text{ dvd}$
 $(\sum_{k=N..n}. (f - g) \$ k * (\text{Rep-prime } p) \wedge \text{nat } (k - N)))$
if $N \leq \min (\text{fls-subdegree } f) (\text{fls-subdegree } g)$
 <proof>

lemma *fls-p-equal0-extended-sum-iff*:

$f \simeq_p 0 \longleftrightarrow$
 $(\forall n. (\text{Rep-prime } p) \wedge \text{Suc } n \text{ dvd}$
 $(\sum_{k \leq n}. f \$ (N + \text{int } k) * (\text{Rep-prime } p) \wedge k))$
if $N \leq \text{fls-subdegree } f$
 <proof>

lemma *fls-p-equal0-extended-intsum-iff*:

$f \simeq_p 0 \longleftrightarrow$
 $(\forall n \geq N. (\text{Rep-prime } p) \wedge \text{Suc } (\text{nat } (n - N)) \text{ dvd}$
 $(\sum_{k=N..n}. f \$ k * (\text{Rep-prime } p) \wedge \text{nat } (k - N)))$
if $N \leq \text{fls-subdegree } f$
 <proof>

end

context

fixes $p :: 'a :: \text{nontrivial-factorial-idom prime}$
begin

lemma *fls-1-not-p-equal-0*: $(1 :: 'a \text{ fls}) \neg \simeq_p 0$
 <proof>

lemma *p-equal-fls0-nonneg-subdegree*:

$f \simeq_p 0 \longleftrightarrow \text{fps-primevanishes } p (\text{fls-regpart } f)$
if $\text{fls-subdegree } f \geq 0$
 <proof>

lemma *p-equal-fls-nonneg-subdegree*:

$f \simeq_p g \longleftrightarrow \text{fps-primevanishes } p (\text{fls-regpart } (f - g))$
if $\text{fls-subdegree } f \geq 0$ **and** $\text{fls-subdegree } g \geq 0$
 <proof>

lemma *p-equal-fls-trans-to-0*:

$f \simeq_p g$ **if** $f \simeq_p 0$ **and** $g \simeq_p 0$ **for** $f \ g :: 'a \text{ fls}$
 <proof>

lemma *p-equal-fls-trans*:

```

     $f \simeq_p g \implies g \simeq_p h \implies$ 
     $f \simeq_p h$ 
    for  $f g h :: 'a \text{ fls}$ 
     $\langle \text{proof} \rangle$ 

end

global-interpretation  $p\text{-equality-fls}$ :
   $p\text{-equality-1}$ 
   $p\text{-equal} :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow$ 
   $'a \text{ fls} \Rightarrow 'a \text{ fls} \Rightarrow \text{bool}$ 
   $\langle \text{proof} \rangle$ 

overloading
   $\text{globally-p-equal-fls} \equiv$ 
   $\text{globally-p-equal} ::$ 
   $'a :: \text{nontrivial-factorial-idom fls} \Rightarrow 'a \text{ fls} \Rightarrow \text{bool}$ 
begin

definition  $\text{globally-p-equal-fls} ::$ 
   $'a :: \text{nontrivial-factorial-idom fls} \Rightarrow 'a \text{ fls} \Rightarrow \text{bool}$ 
  where  $\text{globally-p-equal-fls}[\text{simp}]$ :
     $\text{globally-p-equal-fls} = p\text{-equality-fls}.\text{globally-p-equal}$ 

end

context
  fixes  $p :: 'a :: \text{nontrivial-factorial-idom prime}$ 
begin

lemma  $p\text{-equal-fls-const-iff}$ :
   $(\text{fls-const } c) \simeq_p (\text{fls-const } d) \longleftrightarrow c = d \text{ for } c d :: 'a$ 
   $\langle \text{proof} \rangle$ 

lemma  $p\text{-equal-fls-const-0-iff}$ :
   $(\text{fls-const } c) \simeq_p 0 \longleftrightarrow c = 0 \text{ for } c :: 'a$ 
   $\langle \text{proof} \rangle$ 

lemma  $p\text{-equal-fls-X-intpow-iff}$ :
   $((\text{fls-X-intpow } m :: 'a \text{ fls}) \simeq_p (\text{fls-X-intpow } n)) = (m = n)$ 
   $\langle \text{proof} \rangle$ 

lemma  $\text{fls-X-intpow-nequiv0}$ :  $(\text{fls-X-intpow } m :: 'a \text{ fls}) \not\simeq_p 0$ 
   $\langle \text{proof} \rangle$ 

end

```

4.4.2 Depth as maximum subdegree of equivalent series

Each nonzero equivalence class will have a maximum subdegree.

lemma *no-greatest-p-equal-subdegree:*

$f \simeq_p 0$
if $\forall n :: \text{int. } (\exists g. \text{fls-subdegree } g > n \wedge f \simeq_p g)$
for $p :: 'a::\text{nontrivial-factorial-comm-ring prime}$ **and** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

abbreviation

p-equal-fls-simplified-set $p \ f \equiv$
 $\{g. f \simeq_p g \wedge \text{fls-subdegree } g \geq \text{fls-subdegree } f\}$
— We restrict to those equivalent series that represent a simplification of the given series so that the image of the collection under taking subdegrees is finite.

context

fixes $p :: 'a::\text{nontrivial-factorial-comm-ring prime}$
and $f :: 'a \text{ fls}$
begin

lemma *p-equal-fls-simplified-nempty:*

$\text{fls-subdegree } ' (p\text{-equal-fls-simplified-set } p \ f) \neq \{\}$
 $\langle \text{proof} \rangle$

lemma *p-equal-fls-simplified-finite:*

$\text{finite } (\text{fls-subdegree } ' (p\text{-equal-fls-simplified-set } p \ f))$ **if** $f \not\simeq_p 0$
 $\langle \text{proof} \rangle$

end

overloading

p-depth-fls $\equiv$
 $p\text{-depth} :: 'a::\text{nontrivial-factorial-comm-ring prime} \Rightarrow 'a \text{ fls} \Rightarrow \text{int}$
p-depth-const $\equiv$
 $p\text{-depth} :: 'a::\text{nontrivial-factorial-comm-ring prime} \Rightarrow 'a \Rightarrow \text{int}$
begin

definition

p-depth-fls $:: 'a::\text{nontrivial-factorial-comm-ring prime} \Rightarrow 'a \text{ fls} \Rightarrow \text{int}$
where
p-depth-fls $p \ f \equiv$
 $(\text{if } f \simeq_p 0 \text{ then } 0 \text{ else } \text{Max } (\text{fls-subdegree } ' (p\text{-equal-fls-simplified-set } p \ f)))$

definition

p-depth-const $:: 'a::\text{nontrivial-factorial-comm-ring prime} \Rightarrow 'a \Rightarrow \text{int}$
where
p-depth-const $p \ c \equiv p\text{-depth } p \ (\text{fls-const } c)$

end

context

fixes $p :: 'a::\text{nontrivial-factorial-comm-ring prime}$
begin

lemma $p\text{-depth-fls-eqI}$:

$f^{\circ p} = \text{fls-subdegree } g$
if $f \neg \simeq_p 0$ **and** $f \simeq_p g$
and $\bigwedge h. f \simeq_p h \implies \text{fls-subdegree } h \leq \text{fls-subdegree } g$
for $f g :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-fls-p-equal0[simp]}$:

$(f^{\circ p}) = 0$ **if** $f \simeq_p 0$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-fls0[simp]}$: $(0 :: 'a \text{ fls})^{\circ p} = 0$

$\langle \text{proof} \rangle$

lemma $p\text{-depth-fls-n0D}$:

$f^{\circ p} = \text{Max } (\text{fls-subdegree } ' (p\text{-equal-fls-simplified-set } p f))$
if $f \neg \simeq_p 0$
for $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-flsE}$:

fixes $f :: 'a \text{ fls}$
assumes $f \neg \simeq_p 0$
obtains g **where** $f \simeq_p g$ **and** $\text{fls-subdegree } g = (f^{\circ p})$
and $\forall h. f \simeq_p h \implies \text{fls-subdegree } h \leq \text{fls-subdegree } g$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-fls-ge-p-equal-subdegree}$:

$(f^{\circ p}) \geq \text{fls-subdegree } g$
if $f \neg \simeq_p 0$ **and** $f \simeq_p g$
for $f g :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-fls-ge-p-equal-subdegree}'$:

$(f^{\circ p}) \geq \text{fls-subdegree } f$ **if** $f \neg \simeq_p 0$ **for** $f g :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma $p\text{-equal-fls-simplified-set-shift}$:

$p\text{-equal-fls-simplified-set } p (\text{fls-shift } n f) = \text{fls-shift } n ' p\text{-equal-fls-simplified-set } p f$
if $f \neg \simeq_p 0$
for $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *p-equal-fls-simplified-set-shift-subdegrees*:
 $\text{fls-subdegree } 'p\text{-equal-fls-simplified-set } p \ (\text{fls-shift } n \ f) =$
 $(\lambda x. x - n) \text{ 'fls-subdegree } 'p\text{-equal-fls-simplified-set } p \ f$
if $f \neg \simeq_p 0$
for $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *p-depth-fls-shift*:
 $(\text{fls-shift } n \ f)^{\circ p} = (f^{\circ p}) - n$ **if** $f \neg \simeq_p 0$
for $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

end

context
fixes $p :: 'a :: \text{nontrivial-factorial-idom prime}$
and $f \ g :: 'a \text{ fls}$
begin

lemma *p-depth-fls-p-equal*: $(f^{\circ p}) = (g^{\circ p})$ **if** $f \simeq_p g$
 $\langle \text{proof} \rangle$

lemma *p-depth-fls-can-simplify-iff*:
 $p \text{ pdvd } g \ \S\ (\text{fls-subdegree } g) = (\text{fls-subdegree } g < (f^{\circ p}))$
if $f\text{-nequiv0}: f \neg \simeq_p 0$
and $fg\text{-equiv}: f \simeq_p g$
 $\langle \text{proof} \rangle$

lemma *p-depth-fls-rep-iff*:
 $(f^{\circ p}) = \text{fls-subdegree } g \longleftrightarrow$
 $\neg p \text{ pdvd } g \ \S\ (\text{fls-subdegree } g)$
if $f \neg \simeq_p 0$ **and** $f \simeq_p g$
 $\langle \text{proof} \rangle$

end

global-interpretation *p-depth-fls*:
 $p\text{-equality-depth-no-zero-divisors-1}$
 $p\text{-equal} :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow$
 $'a \text{ fls} \Rightarrow 'a \text{ fls} \Rightarrow \text{bool}$
 $p\text{-depth} :: 'a \text{ prime} \Rightarrow 'a \text{ fls} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

context
fixes $p :: 'a :: \text{nontrivial-factorial-idom prime}$
begin

lemma *p-depth-fls1[simp]*: $(1 :: 'a \text{ fls})^{\circ p} = 0$
 $\langle \text{proof} \rangle$

lemma *p-depth-fls-X*: $(fls-X::'a\ fls)^{\circ p} = 1$
 $\langle proof \rangle$

lemma *p-depth-fls-X-inv*: $(fls-X-inv::'a\ fls)^{\circ p} = -1$
 $\langle proof \rangle$

lemma *p-depth-fls-X-intpow*: $((fls-X-intpow\ n)::'a\ fls)^{\circ p} = n$
 $\langle proof \rangle$

lemma *p-depth-const*:
 $c^{\circ p} = pmultiplicity\ p\ c$ **for** $c :: 'a$
 $\langle proof \rangle$

lemma *p-depth-prime-const*: $(Rep-prime\ p)^{\circ p} = 1$
 $\langle proof \rangle$

end

4.5 Reduction relative to a prime

4.5.1 Of scalars

class *nontrivial-euclidean-ring-cancel* = *nontrivial-factorial-comm-ring* + *euclidean-semiring-cancel*
begin
subclass *nontrivial-factorial-idom* $\langle proof \rangle$
subclass *euclidean-ring-cancel* $\langle proof \rangle$
end

instance *int* :: *nontrivial-euclidean-ring-cancel* $\langle proof \rangle$

abbreviation *pdiv* ::
 $'a \Rightarrow 'a::nontrivial-euclidean-ring-cancel\ prime \Rightarrow 'a$ (**infix** *pdiv* 70)
where $a\ pdiv\ p \equiv a\ div\ (Rep-prime\ p)$

consts *p-modulo* :: $'a \Rightarrow 'b\ prime \Rightarrow 'a$ (**infixl** *pmod* 70)

overloading *p-modulo-scalar* $\equiv p-modulo$:: $'a \Rightarrow 'a\ prime \Rightarrow 'a$
begin
definition *p-modulo-scalar* ::
 $'a::nontrivial-euclidean-ring-cancel \Rightarrow 'a\ prime \Rightarrow 'a$
where *p-modulo-scalar-def[simp]*: $p-modulo-scalar\ a\ p \equiv a\ mod\ (Rep-prime\ p)$
end

class *nontrivial-unique-euclidean-ring* = *nontrivial-factorial-comm-ring* + *unique-euclidean-semiring*
+ **assumes** *division-segment-prime*: $prime\ p \Longrightarrow division-segment\ p = 1$
begin
subclass *nontrivial-euclidean-ring-cancel* $\langle proof \rangle$
end

instance *int* :: *nontrivial-unique-euclidean-ring*
 <proof>

lemma *division-segment-1*: *division-segment* (*1*::'*a*::*unique-euclidean-semiring*) =
 1
 <proof>

lemma *one-pdiv*: (*1*::'*a*) *pdiv* *p* = 0
 and *one-pmod*: (*1*::'*a*) *pmod* *p* = 1
 for *p* :: '*a*::*nontrivial-unique-euclidean-ring prime*
 <proof>

4.5.2 Inductive partial reduction of formal Laurent series

This inductive function reduces one term at a time. The *nat* argument specifies how far beyond the subdegree to reduce, so that all terms to *one less* than the *nat* argument beyond the subdegree are reduced, and the term at the *nat* argument beyond the subdegree is updated to include the carry.

primrec *fls-partially-p-modulo-rep* ::
 '*a*::*nontrivial-euclidean-ring-cancel fls* $\Rightarrow$
nat $\Rightarrow$ '*a prime* $\Rightarrow$ (*int* $\Rightarrow$ '*a*)
where *fls-partially-p-modulo-rep* *f* 0 *p* = (\$\$) *f* |
fls-partially-p-modulo-rep *f* (*Suc* *N*) *p* =
 (*fls-partially-p-modulo-rep* *f* *N* *p*)(
 fls-subdegree *f* + *int* *N* :=
 (*fls-partially-p-modulo-rep* *f* *N* *p* (*fls-subdegree* *f* + *int* *N*)) *pmod* *p*,
 fls-subdegree *f* + *int* *N* + 1 :=
 f \$\$ (*fls-subdegree* *f* + *int* *N* + 1) +
 (*fls-partially-p-modulo-rep* *f* *N* *p* (*fls-subdegree* *f* + *int* *N*)) *pdiv* *p*
)

lemma *fls-partially-p-modulo-rep-zero[simp]*:
fls-partially-p-modulo-rep 0 *N* *p* *n* = 0
 <proof>

lemma *fls-partially-p-modulo-rep-zeros-below*:
fls-partially-p-modulo-rep *f* *N* *p* *n* = 0 **if** *n* < *fls-subdegree* *f*
 <proof>

lemma *fls-partially-p-modulo-rep-func-lower-bound*:
 $\forall n < \text{fls-subdegree } f. \text{ fls-partially-p-modulo-rep } f \text{ (nat } (n + 1 - \text{fls-subdegree } f))$
p *n* = 0
 <proof>

lemma *fls-partially-p-modulo-rep-agrees*:
fls-partially-p-modulo-rep *f* *N* *p* *n* = *fls-partially-p-modulo-rep* *f* *M* *p* *n*
if *n* $\leq$ *fls-subdegree* *f* + *int* *N* - 1 **and** *M* $\geq$ *N*
 <proof>

lemma *fls-partially-p-modulo-rep-func-partially*:

$n > \text{fls-subdegree } f + \text{int } N \implies \text{fls-partially-p-modulo-rep } f \text{ } N \text{ } p \text{ } n = f \text{ } \$\$ \text{ } n$
 $\langle \text{proof} \rangle$

lemma *fls-partially-p-modulo-rep-cong*:

$\text{fls-partially-p-modulo-rep } f \text{ } N \text{ } p \text{ } n = \text{fls-partially-p-modulo-rep } g \text{ } N \text{ } p \text{ } n$
if $\text{fls-subdegree } f = \text{fls-subdegree } g$ **and** $\forall k \leq n. f \text{ } \$\$ \text{ } k = g \text{ } \$\$ \text{ } k$
 $\langle \text{proof} \rangle$

lemma *fls-partially-reduced-nth-vs-rep*:

$N \leq n + 1 - \text{fls-subdegree } f \implies \text{fls-partially-p-modulo-rep } f \text{ } N \text{ } p = (\$\$) \text{ } f$
if $\forall k \leq n. (f \text{ } \$\$ \text{ } k) \text{ } p \text{ div } p = 0$
 $\langle \text{proof} \rangle$

lemma *fls-partially-p-modulo-rep-carry*:

$(\sum_{k=\text{fls-subdegree } f}^N. (f \text{ } \$\$ \text{ } k - \text{fls-partially-p-modulo-rep } f \text{ } (\text{nat } (N - \text{fls-subdegree } f)) \text{ } p \text{ } k) * (\text{Rep-prime } p) \wedge \text{nat } (k - \text{fls-subdegree } f)) = 0$
 $\langle \text{proof} \rangle$

4.5.3 Full reduction of formal Laurent series

Now we collect the reduced terms into a formal Laurent series.

overloading $p\text{-modulo-fls} \equiv p\text{-modulo} :: 'a \text{ fls} \Rightarrow 'a \text{ prime} \Rightarrow 'a \text{ fls}$
begin

definition $p\text{-modulo-fls} ::$

$'a :: \text{nontrivial-euclidean-ring-cancel } \text{fls} \Rightarrow 'a \text{ prime} \Rightarrow 'a \text{ fls}$

where

$p\text{-modulo-fls } f \text{ } p \equiv$

$\text{Abs-fls } (\lambda n. \text{fls-partially-p-modulo-rep } f \text{ } (\text{nat } (n + 1 - \text{fls-subdegree } f)) \text{ } p \text{ } n)$

end

lemma $p\text{-modulo-fls-nth}$:

$(f \text{ } p \text{ mod } p) \text{ } \$\$ \text{ } n = \text{fls-partially-p-modulo-rep } f \text{ } (\text{nat } (n + 1 - \text{fls-subdegree } f)) \text{ } p \text{ } n$
 $\langle \text{proof} \rangle$

lemma $p\text{-modulo-fls-nth}'$:

$\text{int } N \geq n + 1 - \text{fls-subdegree } f \implies$
 $(f \text{ } p \text{ mod } p) \text{ } \$\$ \text{ } n = \text{fls-partially-p-modulo-rep } f \text{ } N \text{ } p \text{ } n$
 $\langle \text{proof} \rangle$

context

fixes $p :: 'a :: \text{nontrivial-euclidean-ring-cancel } \text{prime}$

begin

lemma $\text{fls-pmod-reduced}[simp]$: $((f \text{ } p \text{ mod } p) \text{ } \$\$ \text{ } n) \text{ } p \text{ div } p = 0$
 $\langle \text{proof} \rangle$

lemma *fls-pmod-pmod-nth[simp]*: $((f \text{ pmod } p) \text{ \&\& } n) \text{ pmod } p = (f \text{ pmod } p) \text{ \&\& } n$ **for**
 $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pmod-carry*:
 $(\sum_{k=\text{fls-subdegree } f..N}. (f - (f \text{ pmod } p)) \text{ \&\& } k * (\text{Rep-prime } p) \wedge \text{nat } (k - \text{fls-subdegree } f)) =$
 $(\text{fls-partially-p-modulo-rep } f (\text{nat } (N + 1 - \text{fls-subdegree } f)) p (N + 1) - f \text{ \&\& } (N + 1)) *$
 $(\text{Rep-prime } p) \wedge \text{nat } (N + 1 - \text{fls-subdegree } f)$
 $\langle \text{proof} \rangle$

lemma *fls-pmod-carry'*:
 $(\sum_{k=\text{fls-subdegree } f..N}. (f - (f \text{ pmod } p)) \text{ \&\& } k * (\text{Rep-prime } p) \wedge \text{nat } (k - \text{fls-subdegree } f)) =$
 $(\text{fls-partially-p-modulo-rep } f (\text{nat } (N - \text{fls-subdegree } f)) p N \text{ pdiv } p) *$
 $(\text{Rep-prime } p) \wedge \text{Suc } (\text{nat } (N - \text{fls-subdegree } f))$
if $N \geq \text{fls-subdegree } f$
 $\langle \text{proof} \rangle$

lemma *fls-p-reduced-is-zero-iff*:
 $f = 0 \longleftrightarrow f \simeq_p 0 \wedge (\forall n. f \text{ \&\& } n \text{ pdiv } p = 0)$
for $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *p-modulo-fls-eq0[simp]*:
 $f \text{ pmod } p = 0$ **if** $f \simeq_p 0$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pmod-eq0-iff*:
 $f \text{ pmod } p = 0 \longleftrightarrow f \simeq_p 0$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pmod-equiv*: $f \simeq_p f \text{ pmod } p$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-p-reduced-subdegree-unique*:
 $\text{fls-subdegree } g = (f^{\circ p})$
if $f \simeq_p g$ **and** $\forall n. (g \text{ \&\& } n) \text{ pdiv } p = 0$
for $f g :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pmod-subdegree*:
 $\text{fls-subdegree } (f \text{ pmod } p) = (f^{\circ p})$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pmod-depth*:
 $(f \text{ pmod } p)^{\circ p} = (f^{\circ p})$ **for** $f :: 'a \text{ fls}$

$\langle \text{proof} \rangle$

lemma *fls-pmod-subdegree-ge*:

$\text{fls-subdegree } (f \text{ pmod } p) \geq \text{fls-subdegree } f \text{ if } f \not\simeq_p 0$

for $f :: 'a \text{ fls}$

$\langle \text{proof} \rangle$

lemma *fls-pmod-nth-below-subdegree*:

$(f \text{ pmod } p) \text{ \textit{\$} \$ } n = 0 \text{ if } n < \text{fls-subdegree } f \text{ for } f :: 'a \text{ fls}$

$\langle \text{proof} \rangle$

lemma *fls-pmod-nth-cong*:

$(f \text{ pmod } p) \text{ \textit{\$} \$ } n = (g \text{ pmod } p) \text{ \textit{\$} \$ } n \text{ if } \forall k \leq n. f \text{ \textit{\$} \$ } k = g \text{ \textit{\$} \$ } k \text{ for } f, g :: 'a \text{ fls}$

$\langle \text{proof} \rangle$

lemma *fls-partially-reduced-nth-vs-pmod*:

$(f \text{ pmod } p) \text{ \textit{\$} \$ } n = f \text{ \textit{\$} \$ } n \text{ if } \forall k \leq n. (f \text{ \textit{\$} \$ } k) \text{ pdiv } p = 0$

$\langle \text{proof} \rangle$

lemma *fls-pmod-eqI*:

$f \text{ pmod } p = g \text{ if } f \simeq_p g \text{ and } \forall n. (g \text{ \textit{\$} \$ } n) \text{ pdiv } p = 0$

$\langle \text{proof} \rangle$

lemma *fls-pmod-eq-pmod-iff*:

$f \text{ pmod } p = g \text{ pmod } p \longleftrightarrow f \simeq_p g \text{ for } f, g :: 'a \text{ fls}$

$\langle \text{proof} \rangle$

lemma *fls-pmod-equiv-pmod-iff*:

$(f \text{ pmod } p) \simeq_p (g \text{ pmod } p) \longleftrightarrow f \simeq_p g$

for $f, g :: 'a \text{ fls}$

$\langle \text{proof} \rangle$

lemma *fls-p-modulo-iff*:

$f \text{ pmod } p = g \longleftrightarrow f \simeq_p g \wedge (\forall n. (g \text{ \textit{\$} \$ } n) \text{ pdiv } p = 0)$

$\langle \text{proof} \rangle$

lemma *fls-pmod-pmod[simp]*: $(f \text{ pmod } p) \text{ pmod } p = f \text{ pmod } p \text{ for } f :: 'a \text{ fls}$

$\langle \text{proof} \rangle$

lemma *fls-pmod-cong*:

$f \text{ pmod } p = g \text{ pmod } p \text{ if } f \simeq_p g \text{ for } f, g :: 'a \text{ fls}$

$\langle \text{proof} \rangle$

end

4.5.4 Algebraic properties of reduction

lemma *fls-pmod-1[simp]*:

$(1 :: 'a \text{ fls}) \text{ pmod } p = 1 \text{ for } p :: 'a :: \text{nontrivial-unique-euclidean-ring prime}$

$\langle proof \rangle$

context

fixes $p :: 'a::nontrivial-euclidean-ring-cancel$ *prime*
begin

lemma *fls-pmod-0*: $(0 :: 'a \text{ fls}) \text{ pmod } p = 0$
 $\langle proof \rangle$

lemma *fls-pmod-uminus-equiv*: $(-f) \text{ pmod } p \simeq_p -(f \text{ pmod } p)$ **for** $f :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pmod-equiv-add*:
 $(f + g) \text{ pmod } p = (f' + g') \text{ pmod } p$ **if** $f \simeq_p f'$ **and** $g \simeq_p g'$
for $f f' g g' :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pmod-const-add*:
 $(\text{fls-const } c + f) \text{ pmod } p = \text{fls-const } c + f \text{ pmod } p$ **if** $\text{fls-subdegree } f > 0$ **and** $c \text{ pdiv } p = 0$
for $c :: 'a$ **and** $f :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pmod-add-equiv*:
 $(f + g) \text{ pmod } p \simeq_p (f \text{ pmod } p) + (g \text{ pmod } p)$ **for** $f g :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pmod-mismatched-add-nth*:
 $n < \text{fls-subdegree } g \implies ((f + g) \text{ pmod } p) \$\$ n = (f \text{ pmod } p) \$\$ n$
 $((f + g) \text{ pmod } p) \$\$ (\text{fls-subdegree } g) =$
 $((f \text{ pmod } p) \$\$ (\text{fls-subdegree } g) + g \$\$ (\text{fls-subdegree } g)) \text{ pmod } p$
if $\text{fls-subdegree } g > \text{fls-subdegree } f$
for $f :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pmod-equiv-minus*:
 $(f - g) \text{ pmod } p = (f' - g') \text{ pmod } p$ **if** $f \simeq_p f'$ **and** $g \simeq_p g'$
for $f f' g g' :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pmod-minus-equiv*:
 $(f - g) \text{ pmod } p \simeq_p ((f \text{ pmod } p) - (g \text{ pmod } p))$ **for** $f g :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pmod-equiv-times*:
 $(f * g) \text{ pmod } p = (f' * g') \text{ pmod } p$ **if** $f \simeq_p f'$ **and** $g \simeq_p g'$
for $f f' g g' :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pmod-times-equiv*:
 $(f * g) \text{ pmod } p \simeq_p (f \text{ pmod } p) * (g \text{ pmod } p)$ **for** $f \text{ g} :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pmod-diff-cancel*:
 $(f \text{ pmod } p) \text{ \&\& } n = (g \text{ pmod } p) \text{ \&\& } n$ **if** $n < ((f - g)^{\circ p})$
for $f \text{ g} :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pmod-diff-cancel-iff*:
 $((f - g)^{\circ p}) > n \longleftrightarrow$
 $(\forall k \leq n. (f \text{ pmod } p) \text{ \&\& } k = (g \text{ pmod } p) \text{ \&\& } k)$
if $f \not\simeq_p g$
for $f \text{ g} :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

end

4.6 Inversion relative to a prime

4.6.1 Of scalars

class *bezout-semiring* = *semiring-gcd* +
assumes *bezout*: $\exists u \text{ v. } u * x + v * y = \text{gcd } x \text{ y}$

class *bezout-ring* = *ring-gcd* +
assumes *bezout*: $\exists u \text{ v. } u * x + v * y = \text{gcd } x \text{ y}$
begin
subclass *bezout-semiring* $\langle \text{proof} \rangle$
subclass *idom* $\langle \text{proof} \rangle$
end

instance *int* :: *bezout-ring* $\langle \text{proof} \rangle$

class *nontrivial-factorial-euclidean-bezout* = *nontrivial-euclidean-ring-cancel* + *bezout-semiring*

begin
subclass *bezout-ring* $\langle \text{proof} \rangle$
end

class *nontrivial-factorial-unique-euclidean-bezout* =
 $\text{nontrivial-unique-euclidean-ring} + \text{bezout-semiring}$
begin
subclass *nontrivial-factorial-euclidean-bezout* $\langle \text{proof} \rangle$
end

instance *int* :: *nontrivial-factorial-unique-euclidean-bezout* $\langle \text{proof} \rangle$

lemma *prime-residue-inverse-ex1*:
 $\exists !x. x \text{ pdiv } p = 0 \wedge (x * a) \text{ pmod } p = 1$ **if** $\neg p \text{ pdiv } a$

for $a :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout}$ **and** $p :: 'a \text{ prime}$
 $\langle \text{proof} \rangle$

consts $\text{pinverse} :: 'a \Rightarrow 'b \text{ prime} \Rightarrow 'a$
 $((-^{1-}) [51, 51] 50)$

overloading $\text{pinverse-scalar} \equiv \text{pinverse} :: 'a \Rightarrow 'a \text{ prime} \Rightarrow 'a$
begin

definition $\text{pinverse-scalar} ::$
 $'a::\text{nontrivial-factorial-unique-euclidean-bezout} \Rightarrow 'a \text{ prime} \Rightarrow 'a$
where
 $\text{pinverse-scalar } a \ p \equiv$
 $(\text{if } p \text{ pdvd } a \text{ then } 0 \text{ else}$
 $(\text{THE } x. x \text{ pdiv } p = 0 \wedge (x * a) \text{ pmod } p = 1))$
end

context
fixes $p :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout prime}$
begin

lemma $\text{pinverse-scalar}: ((a^{-1p}) * a) \text{ pmod } p = 1$ **if** $\neg p \text{ pdvd } a$ **for** $a :: 'a$
 $\langle \text{proof} \rangle$

lemma $\text{pinverse-scalar-reduced}: (a^{-1p}) \text{ pdiv } p = 0$ **for** $a :: 'a$
 $\langle \text{proof} \rangle$

lemma $\text{pinverse-scalar-nzero}: (a^{-1p}) \neq 0$ **if** $\neg p \text{ pdvd } a$ **for** $a :: 'a$
 $\langle \text{proof} \rangle$

lemma $\text{pinverse-scalar-zero[simp]}: ((0::'a)^{-1p}) = 0$
 $\langle \text{proof} \rangle$

lemma $\text{pinverse-scalar-one[simp]}: ((1::'a)^{-1p}) = 1$
 $\langle \text{proof} \rangle$

end

4.6.2 Inductive partial inversion of formal Laurent series

Inductively define a depth-zero partial inverse representative function. It is effectively a polynomial of degree equal to the *nat* argument.

primrec $\text{fls-partially-pinverse-rep} ::$
 $'a::\text{nontrivial-factorial-unique-euclidean-bezout fls} \Rightarrow \text{nat} \Rightarrow$
 $'a \text{ prime} \Rightarrow (\text{int} \Rightarrow 'a)$
where
 $\text{fls-partially-pinverse-rep } f \ 0 \ p \ n =$
 $(\text{if } n = 0$
 $\text{then } ((f \text{ pmod } p) \ \$\$ (f^{\circ p}))^{-1p}$
 $\text{else } 0)$

```

| fls-partially-pinverse-rep f (Suc N) p =
  (fls-partially-pinverse-rep f N p)(
    int (Suc N) := -(
      (
        (
          fls-shift (fop) (Abs-fls (fls-partially-pinverse-rep f N p)) *
          (f pmod p)
        ) pmod p
      ) $$ (int (Suc N)) * (((f pmod p) $$ (fop))-1p)
    ) pmod p
  )

```

lemma *fls-partially-pinverse-rep-func-lower-bound:*
 $\forall n < 0. \text{fls-partially-pinverse-rep } f \text{ (nat } n) \text{ } p \text{ } n = 0$
<proof>

lemma *fls-partially-pinverse-rep-at-0:*
 $\text{fls-partially-pinverse-rep } f \text{ } N \text{ } p \text{ } 0 = (((f \text{ pmod } p) \text{ } \$\$ (f^{\text{op}}))^{\text{-1p}})$
<proof>

lemma *fls-partially-pinverse-rep-func-lower-bound':*
 $\forall n < 0. \text{fls-partially-pinverse-rep } f \text{ } N \text{ } p \text{ } n = 0$
<proof>

lemmas *Abs-fls-partially-pinverse-rep-nth =*
nth-Abs-fls-lower-bound[OF fls-partially-pinverse-rep-func-lower-bound']

lemma *fls-partially-pinverse-rep-func-truncates:*
 $n > \text{int } N \implies \text{fls-partially-pinverse-rep } f \text{ } N \text{ } p \text{ } n = 0$
<proof>

context
fixes $p :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime}$
begin

lemma *pinverse-scalar-fls-pmod-base:*
 (
 (((f pmod p) \$\$ (f^{op}))^{-1p}) * (f pmod p) \$\$ (f^{op})
) pmod p = 1
if $f \not\simeq_p 0$
for $f :: 'a \text{ fls}$
<proof>

lemma *fls-partially-pinverse-rep-eq-zero:*
 $\text{fls-partially-pinverse-rep } f \text{ } N \text{ } p = 0 \text{ if } f \simeq_p 0$
for $f :: 'a \text{ fls}$
<proof>

lemma *fls-partially-pinverse-rep-nzero-at-0:*

fls-partially-pinverse-rep f N p $0 \neq 0$ **if** $f \neg \simeq_p 0$ **for** $f :: 'a$ *fls*
 ⟨proof⟩

lemma *fls-partially-pinverse-rep-eq-zero-iff*:
fls-partially-pinverse-rep f N $p = 0 \longleftrightarrow f \simeq_p 0$
for $f :: 'a$ *fls*
 ⟨proof⟩

lemma *Abs-fls-partially-pinverse-rep-nzero*:
Abs-fls (*fls-partially-pinverse-rep* f N p) $\neq 0$ **if** $f \neg \simeq_p 0$
for $f :: 'a$ *fls*
 ⟨proof⟩

lemma *Abs-fls-partially-pinverse-rep-eq-zero-iff*:
Abs-fls (*fls-partially-pinverse-rep* f N p) $= 0 \longleftrightarrow f \simeq_p 0$
for $f :: 'a$ *fls*
 ⟨proof⟩

lemma *Abs-fls-partially-pinverse-rep-subdegree*:
fls-subdegree (*Abs-fls* (*fls-partially-pinverse-rep* f N p)) $= 0$ **if** $f \neg \simeq_p 0$
for $f :: 'a$ *fls*
 ⟨proof⟩

lemma *fls-shift-pinv-rep-subdegree*:
fls-subdegree (*fls-shift* m (*Abs-fls* (*fls-partially-pinverse-rep* f N p))) $= -m$
if $f \neg \simeq_p 0$
for $f :: 'a$ *fls*
 ⟨proof⟩

lemma *fls-partially-pinverse-rep-cong*:
fls-partially-pinverse-rep f N p $(n - K) = \text{fls-partially-pinverse-rep } g$ N p $(n - K)$
if $f^{\circ p} = K$ **and** $g^{\circ p} = K$ **and** $\forall k \leq n. f \text{ $$$ } k = g \text{ $$$ } k$
for $f \text{ } g :: 'a$ *fls*
 ⟨proof⟩

lemma *fls-pinv-rep-times-f-subdegree*:
fls-subdegree (*fls-shift* ($f^{\circ p}$) (*Abs-fls* (*fls-partially-pinverse-rep* f N p)) *
 ($f \text{ pmod } p$)
) $= 0$
if $f \neg \simeq_p 0$
for $f :: 'a$ *fls*
 ⟨proof⟩

lemma *Abs-fls-pinv-rep-Suc*:
fixes $f :: 'a$ *fls* **and** $N :: \text{nat}$
defines
 $c \equiv$

```

      -(
        (
          (
            fls-shift (fop) (Abs-fls (fls-partially-pinverse-rep f N p)) *
            (f pmod p)
          ) pmod p
        ) $$ (int (Suc N)) * (((f pmod p) $$ (fop))-1p)
      ) pmod p
assumes f: f  $\neg \simeq_p 0$ 
shows
  Abs-fls (fls-partially-pinverse-rep f (Suc N) p) =
  Abs-fls (fls-partially-pinverse-rep f N p) + fls-shift (- int (Suc N)) (fls-const
c)
<proof>

```

```

lemma fls-pinverse-rep-times-f-nth:
  ((
    fls-shift (fop) (Abs-fls (fls-partially-pinverse-rep f N p)) *
    (f pmod p)
  ) pmod p) $$ n = 0
if f: f  $\neg \simeq_p 0$  and n: n ≥ 1 n ≤ int N
for f :: 'a fls and n :: int
<proof>

```

end

4.6.3 Full inversion of formal Laurent series

Now we collect the terms of the infinite sequence of partial inversion polynomials into a formal Laurent series, shifted to the opposite depth as the original.

overloading *pinverse-fls* $\equiv$ *pinverse* :: 'a fls $\Rightarrow$ 'a prime $\Rightarrow$ 'a fls

begin

definition *pinverse-fls* ::

'a::nontrivial-factorial-unique-euclidean-bezout fls $\Rightarrow$ 'a prime $\Rightarrow$ 'a fls

where

```

pinverse-fls f p  $\equiv$ 
  (if f  $\simeq_p 0$  then 0 else
   fls-shift (fop)
   (Abs-fls ( $\lambda n$ . fls-partially-pinverse-rep f (nat n) p n)))

```

end

context

fixes p :: 'a::nontrivial-factorial-unique-euclidean-bezout prime

begin

lemma *pinverse-fls-nth*:

```

(f-1p) $$ n =
  fls-partially-pinverse-rep f (nat (n + (fop))) p (n + (fop))

```

if $f \neg \simeq_p 0$
 $\langle \text{proof} \rangle$

lemma *pinverse-fls-nth'*:
 $\text{int } N \geq n + (f^{\circ p}) \implies$
 $(f^{-1p}) \text{ \textit{\$ \$ } } n = \textit{fls-partially-pinverse-rep } f \ N \ p \ (n + (f^{\circ p}))$
if $f \neg \simeq_p 0$
 $\langle \text{proof} \rangle$

lemma *pinverse-fls-eq0[simp]*:
 $(f^{-1p}) = 0 \text{ if } f \simeq_p 0 \text{ for } f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-piniv-eq0-iff*:
 $(f^{-1p}) = 0 \iff f \simeq_p 0 \text{ for } f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-piniv-subdegree*:
 $\textit{fls-subdegree } (f^{-1p}) = -(f^{\circ p}) \text{ for } f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-piniv-reduced[simp]*: $((f^{-1p}) \text{ \textit{\$ \$ } } n) \text{ pdiv } p = 0$
 $\langle \text{proof} \rangle$

lemma *pinverse-times-p-modulo-fls*:
 $((f^{-1p}) * (f \text{ pmod } p)) \text{ pmod } p = 1 \text{ if } f \neg \simeq_p 0$
for $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *pinverse-fls*:
 $((f^{-1p}) * f) \text{ pmod } p = 1 \text{ if } f \neg \simeq_p 0 \text{ for } f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *pinverse-fls-mult-equiv1*:
 $((f^{-1p}) * f) \simeq_p 1 \text{ if } f \neg \simeq_p 0 \text{ for } f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *pinverse-fls-mult-equiv1'*:
 $(f * (f^{-1p})) \simeq_p 1 \text{ if } f \neg \simeq_p 0 \text{ for } f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *p-modulo-pinverse-fls[simp]*:
 $(f^{-1p}) \text{ pmod } p = (f^{-1p}) \text{ for } f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-piniv-equiv0-iff*:
 $(f^{-1p}) \simeq_p 0 \iff f \simeq_p 0$
for $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pinv-depth*:

$(f^{-1p})^{\circ p} = -(f^{\circ p})$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pinv-eqI*:

$(f^{-1p}) = g$ **if** $g * f \simeq_p 1$ **and** $\forall n. (g \text{ $$$ } n) \text{ pdiv } p = 0$
for $f g :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pinv-cong*:

$(f^{-1p}) = (g^{-1p})$ **if** $f \simeq_p g$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *pinverse-p-modulo-fls*:

$(f \text{ pmod } p)^{-1p} = (f^{-1p})$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pinv-X-intpow*: $(\text{fls-X-intpow } n :: 'a \text{ fls})^{-1p} = \text{fls-X-intpow } (-n)$
 $\langle \text{proof} \rangle$

lemma *fls-pinv-uminus*: $(-f)^{-1p} = (- (f^{-1p})) \text{ pmod } p$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

end

4.6.4 Algebraic properties of inversion

context

fixes $p :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime}$
begin

lemma *fls-pinv0*: $(0 :: 'a \text{ fls})^{-1p} = 0$
 $\langle \text{proof} \rangle$

lemma *fls-pinv1[simp]*: $(1 :: 'a \text{ fls})^{-1p} = 1$
 $\langle \text{proof} \rangle$

lemma *fls-pinv-pinv*:

$((f^{-1p})^{-1p}) = f \text{ pmod } p$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pinv-pinv-equiv*:

$((f^{-1p})^{-1p}) \simeq_p f$ **for** $f :: 'a \text{ fls}$
 $\langle \text{proof} \rangle$

lemma *fls-pinv-mult*:

$(f * g)^{-1p} = ((f^{-1p}) * (g^{-1p})) \text{ pmod } p$
for $f g :: 'a \text{ fls}$

$\langle proof \rangle$

lemma *fls-pinv-mult-equiv*:
 $((f * g)^{-1p}) \simeq_p (f^{-1p}) * (g^{-1p})$
for $f g :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pinv-pow*:
 $(f \wedge n)^{-1p} = ((f^{-1p}) \wedge n) \text{ pmod } p$
for $f g :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pinv-pow-equiv*:
 $((f \wedge n)^{-1p}) \simeq_p (f^{-1p}) \wedge n \text{ for } f :: 'a \text{ fls}$
 $\langle proof \rangle$

lemma *fls-pinv-diff-cancel-lead-coeff*:
 $((f^{-1p}) - (g^{-1p}))^{op} > -n$
if $f : f \neg\simeq_p 0 \text{ } f^{op} = n$
and $g : g \neg\simeq_p 0 \text{ } g^{op} = n$
and $fg : f \neg\simeq_p g \text{ } ((f - g)^{op}) > n$
for $f g :: 'a \text{ fls}$
 $\langle proof \rangle$

end

4.7 Topology by place relative to indexed depth for formal Laurent series

4.7.1 General pattern for constructing sequence limits

definition *fls-condition-lim* ::
 $'a :: \text{nontrivial-euclidean-ring-cancel prime} \Rightarrow (\text{nat} \Rightarrow 'a \text{ fls}) \Rightarrow$
 $(\text{nat} \Rightarrow \text{nat} \Rightarrow \text{bool}) \Rightarrow 'a \text{ fls}$
where
 $\text{fls-condition-lim } p \text{ } F \text{ } P =$
 $\text{Abs-fls } (\lambda n. (F \text{ } (\text{LEAST } K. P \text{ } (\text{nat } n) \text{ } K) \text{ pmod } p) \text{ } \$\$ \text{ } n)$

lemma *fls-condition-lim-fun-ex-lower-bound*:
fixes $p :: 'a :: \text{nontrivial-euclidean-ring-cancel prime}$
and $F :: \text{nat} \Rightarrow 'a \text{ fls}$
and $P :: \text{int} \Rightarrow \text{nat} \Rightarrow \text{bool}$
defines
 $f \equiv (\lambda n. (F \text{ } (\text{LEAST } K. P \text{ } (\text{nat } n) \text{ } K) \text{ pmod } p) \text{ } \$\$ \text{ } n)$
and
 $N \equiv \min 0 \text{ } (\text{fls-subdegree } ((F \text{ } (\text{LEAST } K. P \text{ } 0 \text{ } K)) \text{ pmod } p))$
shows $\forall n < N. f \text{ } n = 0$
 $\langle proof \rangle$

lemma *fls-condition-lim-nth*:

fls-condition-lim $p F P \text{ } \$\$ n = ((F (LEAST K. P (nat n) K)) \text{ } pmod p) \text{ } \$\$ n$
 <proof>

lemma *fls-pmod-condition-lim*: $(\text{fls-condition-lim } p F P) \text{ } pmod p = \text{fls-condition-lim } p F P$
 <proof>

4.7.2 Completeness

abbreviation *fls-p-limseq-condition* $\equiv p\text{-depth-fls.p-limseq-condition}$
abbreviation *fls-p-limseq* $\equiv p\text{-depth-fls.p-limseq}$
abbreviation *fls-p-cauchy-condition* $\equiv p\text{-depth-fls.p-cauchy-condition}$
abbreviation *fls-p-cauchy* $\equiv p\text{-depth-fls.p-cauchy}$

lemma *fls-p-cauchy-condition-pmod-uniformity*:
 $((F k) \text{ } pmod p) \text{ } \$\$ m = ((F k') \text{ } pmod p) \text{ } \$\$ m$
if *fls-p-cauchy-condition* $p F n K$ **and** $k \geq K$ **and** $k' \geq K$ **and** $m \leq n$
for $p :: 'a::\text{nontrivial-euclidean-ring-cancel prime}$ **and** $F :: \text{nat} \Rightarrow 'a \text{ fls}$
 <proof>

abbreviation
fls-p-cauchy-lim $p X \equiv$
fls-condition-lim $p X (\lambda n. \text{fls-p-cauchy-condition } p X (\text{int } n))$

lemma *fls-p-cauchy-limseq*: *fls-p-limseq* $p F (\text{fls-p-cauchy-lim } p F)$ **if** *fls-p-cauchy* $p F$
 <proof>

lemma *fls-p-cauchy-lim-unique*:
fls-p-cauchy-lim $p F = \text{fls-p-cauchy-lim } p G$
if *fls-p-cauchy* $p G$ **and** $\forall n \geq N. (F n) \simeq_p (G n)$
 <proof>

4.8 Transfer to prime-indexed sequences of formal Laurent series

4.8.1 Equivalence and depth

The equivalence is by divisibility of the difference of each pair of corresponding terms by the index for those terms.

type-synonym $'a \text{ fls-pseq} = 'a \text{ prime} \Rightarrow 'a \text{ fls}$

overloading
p-equal-fls-pseq $\equiv$
p-equal $:: 'a::\text{nontrivial-factorial-semiring prime} \Rightarrow$
 $'a \text{ fls-pseq} \Rightarrow 'a \text{ fls-pseq} \Rightarrow \text{bool}$
p-restrict-fls-pseq $\equiv$
p-restrict $::$
 $'a \text{ fls-pseq} \Rightarrow ('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow 'a \text{ fls-pseq}$

```

p-depth-fls-pseq ≡
  p-depth :: 'a prime ⇒ 'a fls-pseq ⇒ int
global-unfrmzr-pows-fls-pseq ≡
  global-unfrmzr-pows :: ('a prime ⇒ int) ⇒ 'a fls-pseq
begin

definition p-equal-fls-pseq ::
  'a::nontrivial-factorial-semiring prime ⇒
  'a fls-pseq ⇒ 'a fls-pseq ⇒ bool
  where p-equal-fls-pseq-def[simp]: p-equal-fls-pseq p X Y ≡ (X p) ≃p (Y p)

definition p-restrict-fls-pseq ::
  'a::nontrivial-factorial-semiring fls-pseq ⇒
  ('a prime ⇒ bool) ⇒ 'a fls-pseq
  where
    p-restrict-fls-pseq X P ≡ (λp::'a prime. if P p then X p else 0)

definition p-depth-fls-pseq ::
  'a::nontrivial-factorial-semiring prime ⇒ 'a fls-pseq ⇒ int
  where p-depth-fls-pseq-def[simp]: p-depth-fls-pseq p X ≡ (X p)p

definition global-unfrmzr-pows-fls-pseq ::
  ('a::nontrivial-factorial-semiring prime ⇒ int) ⇒ 'a fls-pseq
  where global-unfrmzr-pows-fls-pseq f ≡ (λp::'a prime. fls-X-intpow (f p))

end

global-interpretation p-equality-depth-fls-pseq:
  global-p-equality-depth-no-zero-divisors-1
  p-equal :: 'a::nontrivial-factorial-idom prime ⇒
  'a fls-pseq ⇒ 'a fls-pseq ⇒ bool
  p-restrict ::
  'a fls-pseq ⇒ ('a prime ⇒ bool) ⇒ 'a fls-pseq
  p-depth :: 'a prime ⇒ 'a fls-pseq ⇒ int
  ⟨proof⟩

overloading
  globally-p-equal-fls-pseq ≡
  globally-p-equal :: 'a::nontrivial-factorial-idom fls-pseq ⇒
  'a fls-pseq ⇒ bool
  p-depth-set-fls-pseq ≡
  p-depth-set ::
  'a::nontrivial-factorial-idom prime ⇒ int ⇒ 'a fls-pseq set
  global-depth-set-fls-pseq ≡
  global-depth-set :: int ⇒ 'a fls-pseq set
begin

definition globally-p-equal-fls-pseq ::
  'a::nontrivial-factorial-idom fls-pseq ⇒ 'a fls-pseq ⇒ bool

```

where *globally-p-equal-fls-pseq-def*[*simp*]:
 $\text{globally-p-equal-fls-pseq} = \text{p-equality-depth-fls-pseq}.\text{globally-p-equal}$

definition *p-depth-set-fls-pseq* ::
 $'a::\text{nontrivial-factorial-idom prime} \Rightarrow \text{int} \Rightarrow 'a \text{ fls-pseq set}$
where *p-depth-set-fls-pseq-def*[*simp*]:
 $\text{p-depth-set-fls-pseq} = \text{p-equality-depth-fls-pseq}.\text{p-depth-set}$

definition *global-depth-set-fls-pseq* ::
 $\text{int} \Rightarrow 'a::\text{nontrivial-factorial-idom fls-pseq set}$
where *global-depth-set-fls-pseq-def*[*simp*]:
 $\text{global-depth-set-fls-pseq} = \text{p-equality-depth-fls-pseq}.\text{global-depth-set}$

end

lemma *fls-pseq-globally-reduced*:
 $X \simeq_{\forall p} (\lambda p. (X \text{ p}) \text{ pmod } p)$
for $X :: 'a::\text{nontrivial-euclidean-ring-cancel fls-pseq}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows0-fls-pseq*:
 $(\text{p } 0 :: 'a::\text{nontrivial-factorial-idom prime} \Rightarrow \text{int}) :: 'a \text{ fls-pseq} = 1$
 $\langle \text{proof} \rangle$

context
fixes $p :: 'a::\text{nontrivial-factorial-idom prime}$
begin

lemma *global-unfrmzr-pows-fls-pseq-nequiv0*:
 $(\text{p } f :: 'a \text{ fls-pseq}) \neg\simeq_p 0$ **for** $f :: 'a \text{ prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-fls-pseq*:
 $(\text{p } f :: 'a \text{ fls-pseq})^{\circ p} = f \text{ p}$ **for** $f :: 'a \text{ prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-fls-pseq-pequiv-iff*:
 $(\text{p } f :: 'a \text{ fls-pseq}) \simeq_p (\text{p } g) \longleftrightarrow f \text{ p} = g \text{ p}$
for $f \text{ g} :: 'a \text{ prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

end

lemma *global-unfrmzr-pows-prod-fls-pseq*:
 $(\text{p } f :: 'a \text{ fls-pseq}) * (\text{p } g) = \text{p } (f + g)$
for $f \text{ g} :: 'a::\text{nontrivial-factorial-idom prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

context

fixes $p :: 'a::\text{nontrivial-factorial-idom prime}$
begin

lemma *global-unfrmzr-pows-fls-pseq-nzero*:
 $(p\ f :: 'a\ \text{fls-pseq}) \negsimeq_p 0$ **for** $f :: 'a\ \text{prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *prod-w-global-unfrmzr-pows-fls-pseq*:
 $(X * p\ f)^{op} = (X^{op}) + f\ p$ **if** $X \negsimeq_p 0$
for $X :: 'a\ \text{fls-pseq}$ **and** $f :: 'a\ \text{prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *normalize-depth-fls-pseq*:
 $(X * p\ (\lambda p::'a\ \text{prime}. -(X^{op})))^{op} = 0$
for $X :: 'a\ \text{fls-pseq}$
 $\langle \text{proof} \rangle$

end

lemma *normalized-depth-fls-pseq-product-additive*:
 $(X * p\ (\lambda p::'a\ \text{prime}. -(X^{op}))) * (Y * p\ (\lambda p::'a\ \text{prime}. -(Y^{op}))) = ((X * Y) * p\ (\lambda p::'a\ \text{prime}. -((X^{op}) + (Y^{op}))))$
for $X\ Y :: 'a::\text{nontrivial-factorial-idom fls-pseq}$
 $\langle \text{proof} \rangle$

context
fixes $p :: 'a::\text{nontrivial-factorial-idom prime}$
begin

lemma *normalized-depth-fls-pseq-product-equiv*:
 $(X * p\ (\lambda p::'a\ \text{prime}. -(X^{op}))) * (Y * p\ (\lambda p::'a\ \text{prime}. -(Y^{op}))) \simeq_p ((X * Y) * p\ (\lambda p::'a\ \text{prime}. -((X * Y)^{op})))$
for $X\ Y :: 'a\ \text{fls-pseq}$
 $\langle \text{proof} \rangle$

lemma *trivial-global-unfrmzr-pows-fls-pseq*:
 $(p\ f :: 'a\ \text{fls-pseq}) \simeq_p 1$ **if** $f\ p = 0$ **for** $f :: 'a\ \text{prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *prod-w-trivial-global-unfrmzr-pows-fls-pseq*:
 $X * p\ f \simeq_p X$ **if** $f\ p = 0$
for $f :: 'a\ \text{prime} \Rightarrow \text{int}$ **and** $X :: 'a\ \text{fls-pseq}$
 $\langle \text{proof} \rangle$

end

lemma *pow-global-unfrmzr-pows-fls-pseq*:

$(\mathfrak{p} \ f :: 'a \text{ fls-pseq}) \wedge n = (\mathfrak{p} \ (\lambda p :: 'a \text{ prime. int } n * f \ p))$
for $f :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-fls-pseq-inv*:
 $(\mathfrak{p} \ (-f) :: 'a \text{ fls-pseq}) * (\mathfrak{p} \ f) = 1$
for $f :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-fls-pseq-decomp*:
 $X = ($
 $\quad (X * \mathfrak{p} \ (\lambda p :: 'a \text{ prime. } -(X^{\circ p})) *$
 $\quad \mathfrak{p} \ (\lambda p :: 'a \text{ prime. } X^{\circ p})$
 $\quad)$
for $X :: 'a :: \text{nontrivial-factorial-idom fls-pseq}$
 $\langle \text{proof} \rangle$

context
fixes $p :: 'a :: \text{nontrivial-factorial-idom prime}$
begin

lemma *global-unfrmzr-pth-fls-pseq*:
 $(\mathfrak{p} \ (1 :: ('a \text{ prime} \Rightarrow \text{int})) \ p :: 'a \text{ fls}) \simeq_p$
 $(\text{fls-const } (\text{Rep-prime } p))$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-fls-pseq*:
 $(\mathfrak{p} \ (1 :: ('a \text{ prime} \Rightarrow \text{int})) :: 'a \text{ fls-pseq}) \simeq_p$
 $(\lambda p :: 'a \text{ prime. fls-const } (\text{Rep-prime } p))$
 $\langle \text{proof} \rangle$

end

lemma *normalize-depth-fls-pseqE*:
fixes $X :: 'a :: \text{nontrivial-factorial-idom fls-pseq}$
obtains $X-0$
where $\forall p :: 'a \text{ prime. } X-0^{\circ p} = 0$
and $X = X-0 * \mathfrak{p} \ (\lambda p :: 'a \text{ prime. } X^{\circ p})$
 $\langle \text{proof} \rangle$

4.8.2 Inversion

abbreviation *global-pinverse-fls-pseq* ::
 $'a :: \text{nontrivial-factorial-comm-ring fls-pseq} \Rightarrow 'a \text{ fls-pseq}$
 $((-^{\circ 1 \vee p}) \ [51] \ 50)$
where $\text{global-pinverse-fls-pseq } X \equiv (\lambda p. (X \ p)^{-1p})$

context
fixes $X :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout fls-pseq}$

begin

lemma *global-pinverse-fls-pseq-eq0[simp]*:
 $(X^{-1\forall p}) = 0$ **if** $X \simeq_{\forall p} 0$
 $\langle proof \rangle$

lemma *global-pinverse-fls-pseq-eq0-iff*:
 $(X^{-1\forall p}) = 0 \longleftrightarrow$
 $X \simeq_{\forall p} 0$
 $\langle proof \rangle$

lemma *global-pinverse-fls-pseq-equiv0-iff*:
 $(X^{-1\forall p}) \simeq_{\forall p} 0 \longleftrightarrow$
 $X \simeq_{\forall p} 0$
 $\langle proof \rangle$

lemma *global-left-pinverse-fls-pseq*:
 $(X^{-1\forall p}) * X \simeq_{\forall p}$
 $(\lambda p. \text{ of-bool } (X \neg\simeq_p 0))$
 $\langle proof \rangle$

lemma *global-right-pinverse-fls-pseq*:
 $X * (X^{-1\forall p}) \simeq_{\forall p}$
 $(\lambda p. \text{ of-bool } (X \neg\simeq_p 0))$
 $\langle proof \rangle$

lemma *global-left-pinverse-fls-pseq'*:
 $(X^{-1\forall p}) * X \simeq_p 1$ **if** $X \neg\simeq_p 0$
for $p :: 'a \text{ prime}$
 $\langle proof \rangle$

lemma *global-right-pinverse-fls-pseq'*:
 $X * (X^{-1\forall p}) \simeq_p 1$ **if** $X \neg\simeq_p 0$
for $p :: 'a \text{ prime}$
 $\langle proof \rangle$

lemma *global-pinverse-pinverse-fls-pseq*:
 $((X^{-1\forall p})^{-1\forall p})$
 $\simeq_{\forall p} X$
 $\langle proof \rangle$

lemma *global-pinverse-mult-fls-pseq*:
 $((X * Y)^{-1\forall p}) \simeq_{\forall p}$
 $(X^{-1\forall p}) * (Y^{-1\forall p})$
 $\langle proof \rangle$

lemma *global-pinverse-pow-fls-pseq*:
 $((X \wedge n)^{-1\forall p}) \simeq_{\forall p}$
 $(X^{-1\forall p}) \wedge n$

$\langle proof \rangle$

lemma *global-pinverse-uminus-fls-pseq*:

$((-X)^{-1\forall p}) \simeq_{\forall p}$
 $- (X^{-1\forall p})$

$\langle proof \rangle$

end

lemma *globally-pinverse-pcong*:

$(X^{-1\forall p}) \simeq_p (Y^{-1\forall p})$

if $X \simeq_p Y$

for $p :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout prime}$ **and** $X\ Y :: 'a$
fls-pseq

$\langle proof \rangle$

lemma *globally-pinverse-cong*:

$(X^{-1\forall p}) \simeq_{\forall p}$
 $(Y^{-1\forall p})$

if $X \simeq_{\forall p} Y$

for $X\ Y :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout fls-pseq}$

$\langle proof \rangle$

lemma *globally-pinverse-global-unfrmzr-pows*:

$(\mathfrak{p}\ f :: 'a\ \text{fls-pseq})^{-1\forall p} = \mathfrak{p}\ (-f)$

for $f :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow \text{int}$

$\langle proof \rangle$

lemma *global-pinverse-fls-pseq0*:

$($
 $(0::'a::\text{nontrivial-factorial-unique-euclidean-bezout fls-pseq})^{-1\forall p}$
 $) = 0$

$\langle proof \rangle$

lemma *global-pinverse-fls-pseq1*:

$($
 $(1::'a::\text{nontrivial-factorial-unique-euclidean-bezout fls-pseq})^{-1\forall p}$
 $) = 1$

$\langle proof \rangle$

lemma *global-pinverse-diff-cancel-lead-coeff*:

$(($
 $(X^{-1\forall p}) -$
 $(Y^{-1\forall p})$

$)^{\circ p}) > -n$

if $X \neg\simeq_p 0$ **and** $X^{\circ p} = n$

and $Y \neg\simeq_p 0$ **and** $Y^{\circ p} = n$

and $X \neg\simeq_p Y$ **and** $((X - Y)^{\circ p}) > n$

for $p :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout prime}$

and $X \ Y :: 'a \text{ fls-pseq}$
 $\langle \text{proof} \rangle$

4.8.3 Topology via global bounded depth

abbreviation $\text{fls-pseq-bdd-global-depth} \equiv p\text{-equality-depth-fls-pseq.bdd-global-depth}$

abbreviation $\text{fls-pseq-bdd-global-dist} \equiv p\text{-equality-depth-fls-pseq.bdd-global-dist}$

abbreviation $\text{fls-pseq-globally-cauchy} \equiv p\text{-equality-depth-fls-pseq.globally-cauchy}$

abbreviation

$\text{fls-pseq-global-cauchy-condition} \equiv p\text{-equality-depth-fls-pseq.global-cauchy-condition}$

lemma $\text{fls-pseq-global-cauchy-condition-pmod-uniformity}$:

$((X \ k \ p) \text{ pmod } p) \ \$\$ \ m = ((X \ k' \ p) \text{ pmod } p) \ \$\$ \ m$

if $\text{fls-pseq-global-cauchy-condition } X \ n \ K$ **and** $k \geq K$ **and** $k' \geq K$ **and** $m \leq \text{int } n$

for $p :: 'a :: \text{nontrivial-euclidean-ring-cancel prime}$ **and** $X :: \text{nat} \Rightarrow 'a \text{ fls-pseq}$

$\langle \text{proof} \rangle$

abbreviation

$\text{fls-pseq-cauchy-lim } X \equiv$

$\lambda p. \text{fls-condition-lim } p \ (\lambda n. X \ n \ p) \ (\text{fls-pseq-global-cauchy-condition } X)$

lemma $\text{fls-pseq-cauchy-condition-lim}$:

$\forall_F k \text{ in sequentially. fls-pseq-bdd-global-dist } (X \ k) \ (\text{fls-pseq-cauchy-lim } X) < e$

if $\text{pos: } e > 0$ **and** $\text{cauchy: fls-pseq-globally-cauchy } X$
 $\langle \text{proof} \rangle$

end

theory $p\text{Adic-Product}$

imports

$\text{FLS-Prime-Equiv-Depth}$

$\text{HOL-Computational-Algebra.Fraction-Field}$

$\text{HOL-Analysis.Elementary-Metric-Spaces}$

begin

5 p-Adic fields as equivalence classes of sequences of formal Laurent series

5.1 Preliminaries

lemma $\text{inj-on-vimage-image-eq}$:

$(f - ' (f - ' B)) \cap A = B$ **if** $\text{inj-on } f \ A$ **and** $B \subseteq A$

$\langle \text{proof} \rangle$

5.2 The type definition as a quotient

quotient-type (overloaded) 'a p-adic-prod
 = 'a::nontrivial-factorial-idom fls-pseq / globally-p-equal
 ⟨proof⟩

lemmas rep-p-adic-prod-inverse = Quotient3-abs-rep[OF Quotient3-p-adic-prod]
 and p-adic-prod-lift-globally-p-equal = Quotient3-rel-abs[OF Quotient3-p-adic-prod]
 and globally-p-equal-fls-pseq-equals-rsp = equals-rsp[OF Quotient3-p-adic-prod]
 and p-adic-prod-eq-iff
 = Quotient3-rel-rep[OF Quotient3-p-adic-prod, symmetric]
 and globally-p-equal-fls-pseq-rep-p-adic-prod-refl
 = Quotient3-rep-refl[OF Quotient3-p-adic-prod]
 and globally-p-equal-fls-pseq-rep-abs-p-adic-prod
 = rep-abs-rsp[OF Quotient3-p-adic-prod] rep-abs-rsp-left[OF Quotient3-p-adic-prod]

lemma p-adic-prod-globally-p-equal-self:
 (rep-p-adic-prod (abs-p-adic-prod r)) $\simeq_p$ r
 ⟨proof⟩

lemma rep-p-adic-prod-p-equal-self: rep-p-adic-prod (abs-p-adic-prod r) $\simeq_p$ r
 for p :: 'a::nontrivial-factorial-idom prime and r :: 'a fls-pseq
 ⟨proof⟩

lemma rep-p-adic-prod-set-inverse: abs-p-adic-prod ' (rep-p-adic-prod ' A) = A
 ⟨proof⟩

lemma p-adic-prod-cases [case-names abs-p-adic-prod, cases type: p-adic-prod]:
 C if $\bigwedge X. x = \text{abs-p-adic-prod } X \implies C$
 ⟨proof⟩

lemmas two-p-adic-prod-cases = p-adic-prod-cases[case-product p-adic-prod-cases]
 and three-p-adic-prod-cases =
 p-adic-prod-cases[case-product p-adic-prod-cases[case-product p-adic-prod-cases]]

lemma p-adic-prod-reduced-cases [case-names abs-p-adic-prod]:
 fixes x :: 'a::nontrivial-euclidean-ring-cancel p-adic-prod
 obtains X where x = abs-p-adic-prod X and $\forall p. (X \ p) \text{ pmod } p = X \ p$
 ⟨proof⟩

lemma p-adic-prod-unique-rep:
 $\exists! X. x = \text{abs-p-adic-prod } X \wedge (\forall p. (X \ p) \text{ pmod } p = X \ p)$
 for x :: 'a::nontrivial-euclidean-ring-cancel p-adic-prod
 ⟨proof⟩

abbreviation

reduced-rep-p-adic-prod x $\equiv$
 (THE X. x = abs-p-adic-prod X $\wedge$ ($\forall p. (X \ p) \text{ pmod } p = X \ p$))

lemma *reduced-rep-p-adic-prod-is-rep*:
 $x = \text{abs-p-adic-prod } (\text{reduced-rep-p-adic-prod } x)$
for $x :: 'a::\text{nontrivial-euclidean-ring-cancel } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *reduced-rep-p-adic-prod-is-reduced*:
 $(\text{reduced-rep-p-adic-prod } x \ p) \text{ pmod } p = \text{reduced-rep-p-adic-prod } x \ p$
for $p :: 'a::\text{nontrivial-euclidean-ring-cancel prime}$ **and** $X :: 'a \text{ p-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *abs-p-adic-prod-inverse*:
 $\text{reduced-rep-p-adic-prod } (\text{abs-p-adic-prod } X) = X$ **if** $\forall p. (X \ p) \text{ pmod } p = X \ p$
for $X :: 'a::\text{nontrivial-euclidean-ring-cancel fls-pseq}$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-seq-cases* [*case-names abs-p-adic-prod*]:
 C **if** $\bigwedge F. X = (\lambda n. \text{abs-p-adic-prod } (F \ n)) \implies C$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-seq-reduced-cases* [*case-names abs-p-adic-prod*]:
fixes $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-euclidean-ring-cancel } p\text{-adic-prod}$
obtains F
where $X = (\lambda n. \text{abs-p-adic-prod } (F \ n))$
and $\forall (n::\text{nat}) (p::'a \text{ prime}). (F \ n \ p) \text{ pmod } p = F \ n \ p$
 $\langle \text{proof} \rangle$

5.3 Algebraic instantiations

The product of p-adic fields form a ring.

instantiation *p-adic-prod* :: (*nontrivial-factorial-idom*) *comm-ring-1*
begin

lift-definition *zero-p-adic-prod* :: '*a* *p-adic-prod* **is** *0*::'*a* *fls-pseq* $\langle \text{proof} \rangle$

lift-definition *one-p-adic-prod* :: '*a* *p-adic-prod* **is** *1*::'*a* *fls-pseq* $\langle \text{proof} \rangle$

lift-definition *plus-p-adic-prod* ::
'*a* *p-adic-prod* $\Rightarrow$ '*a* *p-adic-prod* $\Rightarrow$ '*a* *p-adic-prod*
is $\lambda X \ Y. X + Y$
 $\langle \text{proof} \rangle$

lift-definition *uminus-p-adic-prod* :: '*a* *p-adic-prod* $\Rightarrow$ '*a* *p-adic-prod*
is $\lambda X. - X$
 $\langle \text{proof} \rangle$

lift-definition *minus-p-adic-prod* ::
'*a* *p-adic-prod* $\Rightarrow$ '*a* *p-adic-prod* $\Rightarrow$ '*a* *p-adic-prod*
is $\lambda X \ Y. X - Y$
 $\langle \text{proof} \rangle$

lift-definition *times-p-adic-prod* ::
'a p-adic-prod $\Rightarrow$ *'a p-adic-prod* $\Rightarrow$ *'a p-adic-prod*
is $\lambda X Y. X * Y$
 $\langle proof \rangle$

instance
 $\langle proof \rangle$

end

lemma *pow-p-adic-prod-abs-eq*:
 $(abs\text{-}p\text{-adic}\text{-}prod\ X) \wedge^n = abs\text{-}p\text{-adic}\text{-}prod\ (X \wedge^n)$
for $X :: 'a::nontrivial\text{-}factorial\text{-}idom\ fls\text{-}pseq$
 $\langle proof \rangle$

We can create inverses at the nonzero places.

instantiation *p-adic-prod* ::
 $(nontrivial\text{-}factorial\text{-}unique\text{-}euclidean\text{-}bezout)\ \{divide\text{-}trivial, inverse\}$
begin

lift-definition *divide-p-adic-prod* ::
'a p-adic-prod $\Rightarrow$ *'a p-adic-prod* $\Rightarrow$ *'a p-adic-prod*
is $\lambda X Y. X * (Y^{-1 \nabla p})$
 $\langle proof \rangle$

lift-definition *inverse-p-adic-prod* :: *'a p-adic-prod* $\Rightarrow$ *'a p-adic-prod*
is *global-pinverse-fls-pseq*
 $\langle proof \rangle$

instance
 $\langle proof \rangle$

end

5.4 Equivalence and depth relative to a prime

overloading

p-equal-p-adic-prod $\equiv$
p-equal :: *'a::nontrivial-factorial-idom prime* $\Rightarrow$ *'a p-adic-prod* $\Rightarrow$
'a p-adic-prod $\Rightarrow$ *bool*
p-restrict-p-adic-prod $\equiv$
p-restrict :: *'a p-adic-prod* $\Rightarrow$
 $(\text{'a prime} \Rightarrow \text{bool}) \Rightarrow$ *'a p-adic-prod*
p-depth-p-adic-prod $\equiv$ *p-depth* :: *'a prime* $\Rightarrow$ *'a p-adic-prod* $\Rightarrow$ *int*
global-unfrmzr-pows-p-adic-prod $\equiv$
global-unfrmzr-pows :: $(\text{'a prime} \Rightarrow \text{int}) \Rightarrow$ *'a p-adic-prod*
begin

```

lift-definition p-equal-p-adic-prod ::
  'a::nontrivial-factorial-idom prime  $\Rightarrow$ 
    'a p-adic-prod  $\Rightarrow$  'a p-adic-prod  $\Rightarrow$  bool
is p-equal
  <proof>

lift-definition p-restrict-p-adic-prod ::
  'a::nontrivial-factorial-idom p-adic-prod  $\Rightarrow$ 
    ('a prime  $\Rightarrow$  bool)  $\Rightarrow$  'a p-adic-prod
is  $\lambda X P p.$  if P p then X p else 0
  <proof>

lift-definition p-depth-p-adic-prod ::
  'a::nontrivial-factorial-idom prime  $\Rightarrow$  'a p-adic-prod  $\Rightarrow$  int
is p-depth
  <proof>

lift-definition global-unfrmzr-pows-p-adic-prod ::
  ('a::nontrivial-factorial-idom prime  $\Rightarrow$  int)  $\Rightarrow$  'a p-adic-prod
is global-unfrmzr-pows <proof>

end

lemma p-equal-p-adic-prod-abs-eq0: (abs-p-adic-prod X  $\simeq_p$  0) = (X  $\simeq_p$  0)
for p :: 'a::nontrivial-factorial-idom prime and X :: 'a fls-pseq
  <proof>

lemma p-equal-p-adic-prod-abs-eq1: (abs-p-adic-prod X  $\simeq_p$  1) = (X  $\simeq_p$  1)
for p :: 'a::nontrivial-factorial-idom prime and X :: 'a fls-pseq
  <proof>

global-interpretation p-equality-p-adic-prod:
  p-equality-no-zero-divisors-1
  p-equal :: 'a::nontrivial-factorial-idom prime  $\Rightarrow$ 
    'a p-adic-prod  $\Rightarrow$  'a p-adic-prod  $\Rightarrow$  bool
  <proof>

overloading
  globally-p-equal-p-adic-prod  $\equiv$ 
  globally-p-equal ::
    'a::nontrivial-factorial-idom p-adic-prod  $\Rightarrow$  'a p-adic-prod  $\Rightarrow$  bool
begin

definition globally-p-equal-p-adic-prod ::
  'a::nontrivial-factorial-idom p-adic-prod  $\Rightarrow$  'a p-adic-prod  $\Rightarrow$  bool
where globally-p-equal-p-adic-prod-def[simp]:
  globally-p-equal-p-adic-prod = p-equality-p-adic-prod.globally-p-equal

end

```

global-interpretation *global-p-depth-p-adic-prod*:

global-p-equal-depth-no-zero-divisors-w-inj-index-consts

p-equal :: 'a::nontrivial-factorial-idom prime $\Rightarrow$
'a p-adic-prod $\Rightarrow$ 'a p-adic-prod $\Rightarrow$ bool

p-restrict ::
'a p-adic-prod $\Rightarrow$ ('a prime $\Rightarrow$ bool) $\Rightarrow$ 'a p-adic-prod

p-depth :: 'a prime $\Rightarrow$ 'a p-adic-prod $\Rightarrow$ int

$\lambda f. \text{abs-p-adic-prod } (\lambda p. \text{fls-const } (f p))$

Rep-prime

$\langle \text{proof} \rangle$

lemma *p-depth-p-adic-prod-diff-abs-eq*:

$((\text{abs-p-adic-prod } X - \text{abs-p-adic-prod } Y)^{\circ p}) = ((X - Y)^{\circ p})$

for *p* :: 'a::nontrivial-factorial-idom prime **and** *X Y* :: 'a fls-pseq

$\langle \text{proof} \rangle$

lemma *p-depth-p-adic-prod-diff-abs-eq'*:

$((\text{abs-p-adic-prod } X - \text{abs-p-adic-prod } Y)^{\circ p}) = ((X p - Y p)^{\circ p})$

for *p* :: 'a::nontrivial-factorial-idom prime **and** *X Y* :: 'a fls-pseq

$\langle \text{proof} \rangle$

overloading

p-depth-set-p-adic-prod $\equiv$

p-depth-set ::
'a::nontrivial-factorial-idom prime $\Rightarrow$ int $\Rightarrow$ 'a p-adic-prod set

global-depth-set-p-adic-prod $\equiv$

global-depth-set :: int $\Rightarrow$ 'a p-adic-prod set

begin

definition *p-depth-set-p-adic-prod* ::
'a::nontrivial-factorial-idom prime $\Rightarrow$ int $\Rightarrow$ 'a p-adic-prod set

where *p-depth-set-p-adic-prod-def[simp]*:

p-depth-set-p-adic-prod = *global-p-depth-p-adic-prod.p-depth-set*

definition *global-depth-set-p-adic-prod* ::
int $\Rightarrow$ 'a::nontrivial-factorial-idom p-adic-prod set

where *global-depth-set-p-adic-prod-def[simp]*:

global-depth-set-p-adic-prod = *global-p-depth-p-adic-prod.global-depth-set*

end

lemma *p-depth-set-p-adic-prod-eq-projection*:

$((\mathcal{P}_{\mathbb{A}_p}^n) :: 'a \text{ p-adic-prod set}) =$
 $\text{abs-p-adic-prod } '(\mathcal{P}_{\mathbb{A}_p}^n)$

for *p* :: 'a::nontrivial-factorial-idom prime

$\langle \text{proof} \rangle$

lemma *global-depth-set-p-adic-prod-eq-projection*:

$((\mathcal{P}_{\forall p}^n) :: 'a \text{ p-adic-prod set}) =$
 $\text{abs-p-adic-prod } '(\mathcal{P}_{\forall p}^n)$
for $p :: 'a :: \text{nontrivial-factorial-idom prime}$
 $\langle \text{proof} \rangle$

context

fixes $x :: 'a :: \text{nontrivial-factorial-idom p-adic-prod}$
begin

lemma $p\text{-adic-prod-p-restrict-depth}$:
 $(x \text{ prestrict } P)^{\circ p} = (\text{if } P \text{ then } x^{\circ p} \text{ else } 0)$
for $P :: 'a \text{ prime} \Rightarrow \text{bool}$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-global-depth-setI}$:
 $x \in \mathcal{P}_{\forall p}^n$
if $\bigwedge p :: 'a \text{ prime. } x \dashv\sim_p 0 \implies (x^{\circ p}) \geq n$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-global-depth-setI'}$:
 $x \in \mathcal{P}_{\forall p}^n$
if $\bigwedge p :: 'a \text{ prime. } (x^{\circ p}) \geq n$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-global-depth-set-p-restrictI}$:
 $p\text{-restrict } x \text{ } P \in \mathcal{P}_{\forall p}^n$
if $\bigwedge p :: 'a \text{ prime. } P \text{ } p \implies x \in \mathcal{P}_{@p}^n$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-p-depth-setI}$:
 $x \in \mathcal{P}_{@p}^n$
if $x \dashv\sim_p 0 \implies (x^{\circ p}) \geq n$
for $p :: 'a \text{ prime}$
 $\langle \text{proof} \rangle$

end

context

fixes $p :: 'a :: \text{nontrivial-euclidean-ring-cancel prime}$
begin

lemma $p\text{-adic-prod-diff-depth-gt-equiv-trans}$:
 $((x - z)^{\circ p}) > n$
if $xy: x \simeq_p y$
and $yz: y \dashv\sim_p z ((y - z)^{\circ p}) > n$
for $x \ y \ z :: 'a \text{ p-adic-prod}$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-diff-depth-gt0-equiv-trans}$:

$((x - z)^{\circ p}) > n$
if $n \geq 0$ **and** $x \simeq_p y$ **and** $((y - z)^{\circ p}) > n$
for $x \ y \ z :: 'a \text{ p-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-diff-depth-gt-trans:*

$((x - z)^{\circ p}) > n$
if $xy: x \neg\simeq_p y ((x - y)^{\circ p}) > n$
and $yz: y \neg\simeq_p z ((y - z)^{\circ p}) > n$
and $xz: x \neg\simeq_p z$
for $x \ y \ z :: 'a \text{ p-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-diff-depth-gt0-trans:*

$((x - z)^{\circ p}) > n$
if $n \geq 0$
and $((x - y)^{\circ p}) > n$
and $((y - z)^{\circ p}) > n$
and $x \neg\simeq_p z$
for $x \ y \ z :: 'a \text{ p-adic-prod}$
 $\langle \text{proof} \rangle$

end

lemma *p-adic-prod-diff-cancel-lead-coeff:*

$((\text{inverse } x - \text{inverse } y)^{\circ p}) > -n$
if $x : x \neg\simeq_p 0 \ x^{\circ p} = n$
and $y : y \neg\simeq_p 0 \ y^{\circ p} = n$
and $xy: x \neg\simeq_p y ((x - y)^{\circ p}) > n$
for $p :: 'a :: \text{nontrivial-factorial-unique-euclidean-bézout prime}$
and $x \ y :: 'a \text{ p-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-p-adic-prod0:*

$(\text{p } (0 :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{int}) :: 'a \text{ p-adic-prod}) = 1$
 $\langle \text{proof} \rangle$

context

fixes $p :: 'a :: \text{nontrivial-factorial-idom prime}$

begin

lemma *global-unfrmzr-pows-p-adic-prod-nequiv0:*

$(\text{p } f :: 'a \text{ p-adic-prod}) \neg\simeq_p 0$ **for** $f :: 'a \text{ prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-p-adic-prod:*

$(\text{p } f :: 'a \text{ p-adic-prod})^{\circ p} = f \ p$ **for** $f :: 'a \text{ prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

```

lemma global-unfrmzr-pows-p-adic-prod-depth-set:
  fixes  $f :: 'a \text{ prime} \Rightarrow \text{int}$ 
  defines  $n \equiv f \, p$ 
  shows  $(\mathfrak{p} \, f :: 'a \text{ p-adic-prod}) \in \mathcal{P}_{\mathbb{Q}_p}^n$ 
   $\langle \text{proof} \rangle$ 

lemma global-unfrmzr-pows-p-adic-prod-pequiv-iff:
   $((\mathfrak{p} \, f :: 'a \text{ p-adic-prod}) \simeq_p (\mathfrak{p} \, g)) = (f \, p = g \, p)$ 
  for  $f \, g :: 'a \text{ prime} \Rightarrow \text{int}$ 
   $\langle \text{proof} \rangle$ 

end

lemma global-unfrmzr-pows-p-adic-prod-global-depth-set:
   $(\mathfrak{p} \, f :: 'a \text{ p-adic-prod}) \in \mathcal{P}_{\mathbb{Q}_p}^n$ 
  if  $\forall p. f \, p \geq n$  for  $f :: 'a::\text{nontrivial-factorial-idom prime} \Rightarrow \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma global-unfrmzr-pows-p-adic-prod-global-depth-set-0:
   $(\mathfrak{p} \, (\text{int} \circ f) :: 'a \text{ p-adic-prod}) \in \mathcal{O}_{\mathbb{Q}_p}$ 
  for  $f :: 'a::\text{nontrivial-factorial-idom prime} \Rightarrow \text{nat}$ 
   $\langle \text{proof} \rangle$ 

lemma global-unfrmzr-pows-prod-p-adic-prod:
   $(\mathfrak{p} \, f :: 'a \text{ p-adic-prod}) * (\mathfrak{p} \, g) = \mathfrak{p} \, (f + g)$ 
  for  $f \, g :: 'a::\text{nontrivial-factorial-idom prime} \Rightarrow \text{int}$ 
   $\langle \text{proof} \rangle$ 

context
  fixes  $p :: 'a::\text{nontrivial-factorial-idom prime}$ 
begin

lemma global-unfrmzr-pows-p-adic-prod-nzero:
   $(\mathfrak{p} \, f :: 'a \text{ p-adic-prod}) \not\simeq_p 0$  for  $f :: 'a \text{ prime} \Rightarrow \text{int}$ 
   $\langle \text{proof} \rangle$ 

lemma prod-w-global-unfrmzr-pows-p-adic-prod:
   $x \not\simeq_p 0 \implies$ 
   $(x * \mathfrak{p} \, f)^{\circ p} = (x^{\circ p}) + f \, p$ 
  for  $f :: 'a \text{ prime} \Rightarrow \text{int}$  and  $x :: 'a \text{ p-adic-prod}$ 
   $\langle \text{proof} \rangle$ 

lemma p-adic-prod-normalize-depth:
   $(x * \mathfrak{p} \, (\lambda p::'a \text{ prime. } -(x^{\circ p})))^{\circ p} = 0$ 
  for  $x :: 'a \text{ p-adic-prod}$ 
   $\langle \text{proof} \rangle$ 

end

```

lemma *p-adic-prod-normalized-depth-product-equiv:*

$(x * \mathfrak{p} (\lambda p :: 'a \text{ prime. } -(x^{\circ p}))) *$
 $(y * \mathfrak{p} (\lambda p :: 'a \text{ prime. } -(y^{\circ p}))) =$
 $((x * y) * \mathfrak{p} (\lambda p :: 'a \text{ prime. } -((x * y)^{\circ p})))$
for $x \ y :: 'a :: \text{nontrivial-factorial-idom } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

context

fixes $p :: 'a :: \text{nontrivial-factorial-idom prime}$
begin

lemma *trivial-global-unfrmzr-pows-p-adic-prod:*

$f \ p = 0 \implies (\mathfrak{p} \ f :: 'a \ p\text{-adic-prod}) \simeq_p 1$
for $f :: 'a \text{ prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *prod-w-trivial-global-unfrmzr-pows-p-adic-prod:*

$f \ p = 0 \implies x * \mathfrak{p} \ f \simeq_p x$
for $f :: 'a \text{ prime} \Rightarrow \text{int}$ **and** $x :: 'a \ p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

end

lemma *global-unfrmzr-pows-p-adic-prod-pow:*

$(\mathfrak{p} \ f :: 'a \ p\text{-adic-prod})^{\wedge} n = (\mathfrak{p} (\lambda p. \text{int } n * f \ p))$
for $f :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-p-adic-prod-inv:*

$(\mathfrak{p} (-f) :: 'a \ p\text{-adic-prod}) * (\mathfrak{p} \ f) = 1$
for $f :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-p-adic-prod-inverse:*

inverse $(\mathfrak{p} \ f :: 'a \ p\text{-adic-prod}) = \mathfrak{p} (-f)$
for $f :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-p-adic-prod-decomp:*

$x = ($
 $(x * \mathfrak{p} (\lambda p :: 'a \text{ prime. } -(x^{\circ p}))) *$
 $\mathfrak{p} (\lambda p :: 'a \text{ prime. } x^{\circ p})$
 $)$
for $p :: 'a :: \text{nontrivial-factorial-idom prime}$ **and** $x :: 'a \ p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-normalize-depthE:*

fixes $x :: 'a :: \text{nontrivial-factorial-idom } p\text{-adic-prod}$
obtains $x \cdot 0$
where $\forall p :: 'a \text{ prime. } x \cdot 0^{\circ p} = 0$

and $x = x-0 * p$ ($\lambda p :: 'a \text{ prime. } x^{\circ p}$)
 $\langle \text{proof} \rangle$

global-interpretation $p\text{-adic-prod-div-inv}$:

$\text{global-p-equal-depth-div-inv}$
 $p\text{-equal} :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow$
 $'a \text{ p-adic-prod} \Rightarrow 'a \text{ p-adic-prod} \Rightarrow \text{bool}$
 $p\text{-depth} :: 'a \text{ prime} \Rightarrow 'a \text{ p-adic-prod} \Rightarrow \text{int}$
 $p\text{-restrict} ::$
 $'a \text{ p-adic-prod} \Rightarrow ('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow 'a \text{ p-adic-prod}$
 $\langle \text{proof} \rangle$

5.5 Embeddings of the base ring, nat , int , and rat

5.5.1 Embedding of the base ring

abbreviation $p\text{-adic-prod-consts} \equiv \text{global-p-depth-p-adic-prod.consts}$

abbreviation $p\text{-adic-prod-const} \equiv \text{global-p-depth-p-adic-prod.const}$

lemma $p\text{-depth-p-adic-prod-consts}$:

$(p\text{-adic-prod-consts } f)^{\circ p} = ((f \text{ } p)^{\circ p})$

for $p :: 'a :: \text{nontrivial-factorial-idom prime}$ **and** $f :: 'a \text{ prime} \Rightarrow 'a$

$\langle \text{proof} \rangle$

lemmas $p\text{-depth-p-adic-prod-const} = p\text{-depth-p-adic-prod-consts}[of \text{ } \lambda p. \text{ } -]$

lemma $p\text{-adic-prod-global-unfrmr}$:

$(p \text{ } (1 :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{int}) :: 'a \text{ p-adic-prod}) =$
 $p\text{-adic-prod-consts Rep-prime}$

$\langle \text{proof} \rangle$

5.5.2 Embedding of the field of fractions of the base ring

global-interpretation $p\text{-adic-prod-embeds}$:

$\text{global-p-equal-div-inv-w-inj-idom-consts}$
 $p\text{-equal} :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow$
 $'a \text{ p-adic-prod} \Rightarrow 'a \text{ p-adic-prod} \Rightarrow \text{bool}$
 $p\text{-restrict} ::$
 $'a \text{ p-adic-prod} \Rightarrow ('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow 'a \text{ p-adic-prod}$
 $\lambda f. \text{ abs-p-adic-prod } (\lambda p. \text{ fls-const } (f \text{ } p))$
 $\langle \text{proof} \rangle$

global-interpretation $p\text{-adic-prod-depth-embeds}$:

$\text{global-p-equal-depth-div-inv-w-inj-index-consts}$
 $p\text{-equal} :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow$
 $'a \text{ p-adic-prod} \Rightarrow 'a \text{ p-adic-prod} \Rightarrow \text{bool}$
 $p\text{-depth} :: 'a \text{ prime} \Rightarrow 'a \text{ p-adic-prod} \Rightarrow \text{int}$
 $p\text{-restrict} ::$
 $'a \text{ p-adic-prod} \Rightarrow ('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow 'a \text{ p-adic-prod}$
 $\lambda f. \text{ abs-p-adic-prod } (\lambda p. \text{ fls-const } (f \text{ } p))$

Rep-prime
 ⟨proof⟩

abbreviation *p-adic-prod-shift-p-depth* $\equiv$ *p-adic-prod-depth-embeds.shift-p-depth*

abbreviation *p-adic-prod-of-fract* $\equiv$ *p-adic-prod-embeds.const-of-fract*

lemma *p-adic-prod-of-fract-depth*:

(*p-adic-prod-of-fract* (*Fraction-Field.Fract* *a b*) :: '*a p-adic-prod*)^{op} =
 (*a*^{op}) - (*b*^{op})

if *a* ≠ 0 **and** *b* ≠ 0

for *p* :: '*a::nontrivial-factorial-unique-euclidean-bezout prime*

and *a b* :: '*a*

⟨proof⟩

Product of all p-adic embeddings of the field of fractions of the base ring.

abbreviation *p-adic-Fracts-prod* ::

'*a::nontrivial-factorial-unique-euclidean-bezout p-adic-prod set*

(*Q_{v p}*)

where *Q_{v p}* $\equiv$ *p-adic-prod-embeds.range-const-of-fract*

5.5.3 Characteristic zero embeddings

class *nontrivial-factorial-idom-char-0* = *nontrivial-factorial-idom* + *semiring-char-0*

begin

subclass *ring-char-0* ⟨proof⟩

end

instance *int* :: *nontrivial-factorial-idom-char-0* ⟨proof⟩

class *nontrivial-factorial-unique-euclidean-bezout-char-0* =

nontrivial-factorial-unique-euclidean-bezout + *nontrivial-factorial-idom-char-0*

instance *int* :: *nontrivial-factorial-unique-euclidean-bezout-char-0* ⟨proof⟩

instance *p-adic-prod* :: (*nontrivial-factorial-idom-char-0*) *ring-char-0*

⟨proof⟩

global-interpretation *p-adic-prod-consts-char-0*:

global-p-equality-w-inj-consts-char-0

p-equal :: '*a::nontrivial-factorial-idom-char-0 prime* $\Rightarrow$

'*a p-adic-prod* $\Rightarrow$ '*a p-adic-prod* $\Rightarrow$ *bool*

p-restrict ::

'*a p-adic-prod* $\Rightarrow$ ('*a prime* $\Rightarrow$ *bool*) $\Rightarrow$ '*a p-adic-prod*

$\lambda f. \text{abs-}p\text{-adic-prod } (\lambda p. \text{fls-const } (f p))$

⟨proof⟩

global-interpretation *p-adic-prod-div-inv-consts-char-0*:

global-p-equal-div-inv-w-inj-consts-char-0

p-equal :: '*a::nontrivial-factorial-unique-euclidean-bezout-char-0 prime* $\Rightarrow$

$'a \text{ p-adic-prod} \Rightarrow 'a \text{ p-adic-prod} \Rightarrow \text{bool}$
 $p\text{-restrict} ::$
 $'a \text{ p-adic-prod} \Rightarrow ('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow 'a \text{ p-adic-prod}$
 $\lambda f. \text{abs-p-adic-prod } (\lambda p. \text{fls-const } (f p))$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-of-nat-depth}$:
 $(\text{of-nat } n :: 'a \text{ p-adic-prod})^{\circ p} = ((\text{of-nat } n :: 'a)^{\circ p})$
for $p :: 'a :: \text{nontrivial-factorial-idom prime}$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-of-int-depth}$:
 $(\text{of-int } z :: 'a \text{ p-adic-prod})^{\circ p} = ((\text{of-int } z :: 'a)^{\circ p})$
for $p :: 'a :: \text{nontrivial-factorial-idom prime}$
 $\langle \text{proof} \rangle$

abbreviation $p\text{-adic-prod-of-rat} \equiv p\text{-adic-prod-div-inv-consts-char-0.const-of-rat}$

lemma $p\text{-adic-prod-of-rat-depth}$:
 $(p\text{-adic-prod-of-rat } (\text{Rat.Fract } a \ b) :: 'a \text{ p-adic-prod})^{\circ p} =$
 $((\text{of-int } a :: 'a)^{\circ p}) - ((\text{of-int } b :: 'a)^{\circ p})$
if $a \neq 0$ **and** $b \neq 0$
for $p :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout-char-0 prime}$
and $a \ b :: \text{int}$
 $\langle \text{proof} \rangle$

5.6 Topologies on products of p-adic fields

5.6.1 By local place

Completeness

abbreviation $p\text{-adic-prod-p-open-nbhd} \equiv \text{global-p-depth-p-adic-prod.p-open-nbhd}$
abbreviation $p\text{-adic-prod-p-open-nbhds} \equiv \text{global-p-depth-p-adic-prod.p-open-nbhds}$
abbreviation
 $p\text{-adic-prod-p-limseq-condition} \equiv \text{global-p-depth-p-adic-prod.p-limseq-condition}$
abbreviation
 $p\text{-adic-prod-p-cauchy-condition} \equiv \text{global-p-depth-p-adic-prod.p-cauchy-condition}$
abbreviation $p\text{-adic-prod-p-limseq} \equiv \text{global-p-depth-p-adic-prod.p-limseq}$
abbreviation $p\text{-adic-prod-p-cauchy} \equiv \text{global-p-depth-p-adic-prod.p-cauchy}$
abbreviation
 $p\text{-adic-prod-p-open-nbhds-limseq} \equiv \text{global-p-depth-p-adic-prod.p-open-nbhds-LIMSEQ}$
abbreviation
 $p\text{-adic-prod-local-p-open-nbhds} \equiv \text{global-p-depth-p-adic-prod.local-p-open-nbhds}$

lemma $\text{abs-p-adic-prod-seq-cases } [\text{case-names } \text{abs-p-adic-prod}]$:
 $C \text{ if } \bigwedge F. X = (\lambda n. \text{abs-p-adic-prod } (F n)) \implies C$
 $\langle \text{proof} \rangle$

lemma $\text{abs-p-adic-prod-p-limseq-condition}$:

$p\text{-adic-prod-p-limseq-condition } p (\lambda n. \text{abs-p-adic-prod } (F \ n)) (\text{abs-p-adic-prod } X)$
 $n \ K =$
 $\text{fls-p-limseq-condition } p (\lambda n. F \ n \ p) (X \ p) \ n \ K$
 $\langle \text{proof} \rangle$

lemma $\text{abs-p-adic-prod-p-limseq}$:
 $p\text{-adic-prod-p-limseq } p (\lambda n. \text{abs-p-adic-prod } (F \ n)) (\text{abs-p-adic-prod } X) =$
 $\text{fls-p-limseq } p (\lambda n. F \ n \ p) (X \ p)$
 $\langle \text{proof} \rangle$

lemma $\text{abs-p-adic-prod-p-cauchy-condition}$:
 $p\text{-adic-prod-p-cauchy-condition } p (\lambda n. \text{abs-p-adic-prod } (F \ n)) \ n \ K =$
 $\text{fls-p-cauchy-condition } p (\lambda n. F \ n \ p) \ n \ K$
 $\langle \text{proof} \rangle$

lemma $\text{abs-p-adic-prod-p-cauchy}$:
 $p\text{-adic-prod-p-cauchy } p (\lambda n. \text{abs-p-adic-prod } (F \ n)) = \text{fls-p-cauchy } p (\lambda n. F \ n \ p)$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-p-restrict-cauchy-lim}$:
 $p\text{-adic-prod-p-limseq } q$
 $(\lambda n. p\text{-restrict } (X \ n) ((=) \ p))$
 $(\text{abs-p-adic-prod } (\lambda q.$
 $\text{if } q = p \text{ then fls-p-cauchy-lim } p (\lambda n. \text{rep-p-adic-prod } (X \ n) \ p) \text{ else } 0$
 $))$
if $p\text{-adic-prod-p-cauchy } p \ X$
 $\langle \text{proof} \rangle$

global-interpretation $p\text{-complete-p-adic-prod}$:
 $p\text{-complete-global-p-equal-depth}$
 $p\text{-equal} :: 'a :: \text{nontrivial-euclidean-ring-cancel prime} \Rightarrow$
 $'a \ p\text{-adic-prod} \Rightarrow 'a \ p\text{-adic-prod} \Rightarrow \text{bool}$
 $p\text{-restrict} ::$
 $'a \ p\text{-adic-prod} \Rightarrow ('a \ \text{prime} \Rightarrow \text{bool}) \Rightarrow 'a \ p\text{-adic-prod}$
 $p\text{-depth} :: 'a \ \text{prime} \Rightarrow 'a \ p\text{-adic-prod} \Rightarrow \text{int}$
 $\langle \text{proof} \rangle$

global-interpretation $p\text{-complete-p-adic-prod-div-inv}$:
 $p\text{-complete-global-p-equal-depth-div-inv}$
 $p\text{-equal} :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow$
 $'a \ p\text{-adic-prod} \Rightarrow 'a \ p\text{-adic-prod} \Rightarrow \text{bool}$
 $p\text{-depth} :: 'a \ \text{prime} \Rightarrow 'a \ p\text{-adic-prod} \Rightarrow \text{int}$
 $p\text{-restrict} ::$
 $'a \ p\text{-adic-prod} \Rightarrow ('a \ \text{prime} \Rightarrow \text{bool}) \Rightarrow 'a \ p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemmas $p\text{-adic-prod-hensel} = p\text{-complete-p-adic-prod-div-inv.hensel}$

5.6.2 By bounded global depth

The metric

instantiation $p\text{-adic-prod} :: (\text{nontrivial-factorial-idom}) \text{ metric-space}$
begin

definition $\text{dist} = \text{global-p-depth-p-adic-prod.bdd-global-dist}$
declare $\text{dist-p-adic-prod-def} \text{ [simp]}$

definition $\text{uniformity} = \text{global-p-depth-p-adic-prod.bdd-global-uniformity}$
declare $\text{uniformity-p-adic-prod-def} \text{ [simp]}$

definition $\text{open-p-adic-prod} :: 'a \text{ p-adic-prod set} \Rightarrow \text{bool}$
where
 $\text{open-p-adic-prod } U = (\forall x \in U.$
 $\text{eventually } (\lambda(x', y). x' = x \longrightarrow y \in U) \text{ uniformity}$
 $)$

instance
 $\langle \text{proof} \rangle$

end

abbreviation $p\text{-adic-prod-global-depth} \equiv \text{global-p-depth-p-adic-prod.bdd-global-depth}$

lemma $\text{bdd-global-depth-p-adic-prod-abs-eq}$:
 $p\text{-adic-prod-global-depth } (\text{abs-p-adic-prod } X) = \text{fls-pseq-bdd-global-depth } X$
 $\langle \text{proof} \rangle$

lemma $\text{dist-p-adic-prod-abs-eq}$:
 $\text{dist } (\text{abs-p-adic-prod } X) (\text{abs-p-adic-prod } Y) = \text{fls-pseq-bdd-global-dist } X Y$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-metric-is-finer}$:
 $\text{generate-topology } p\text{-adic-prod-p-open-nbhds } U \Longrightarrow \text{open } U$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-metric-is-finer-than-purely-local}$:
 $\text{generate-topology } (p\text{-adic-prod-local-p-open-nbhds } p) U \Longrightarrow$
 $\text{open } U$
for $p :: 'a :: \text{nontrivial-factorial-idom prime}$
and $U :: 'a \text{ p-adic-prod set}$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-limseq-abs-eq}$:
 $(\lambda n. \text{abs-p-adic-prod } (X n)) \longrightarrow \text{abs-p-adic-prod } x$
 $\longleftrightarrow$
 $(\forall e > 0. \forall_F n \text{ in sequentially. fls-pseq-bdd-global-dist } (X n) x < e)$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-bdd-metric-LIMSEQ*:

$X \longrightarrow x = \text{global-p-depth-p-adic-prod.metric-space-by-bdd-depth.LIMSEQ } X \ x$
for $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$ **and** $x :: 'a \text{ } p\text{-adic-prod}$
 ⟨proof⟩

lemma *p-adic-prod-global-depth-set-lim-closed*:

$x \in \mathcal{P}_{\forall p}^n$
if $\text{lim} : X \longrightarrow x$
and $\text{depth}: \forall_F k \text{ in sequentially. } X \ k \in \mathcal{P}_{\forall p}^n$
for $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-factorial-unique-euclidean-bezout } p\text{-adic-prod}$
and $x :: 'a \text{ } p\text{-adic-prod}$
 ⟨proof⟩

lemma *p-adic-prod-metric-ball-multicentre*:

$\text{ball } y \ e = \text{ball } x \ e$ **if** $y \in \text{ball } x \ e$
for $x \ y :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$
 ⟨proof⟩

lemma *p-adic-prod-metric-ball-at-0*:

$\text{ball } (0::'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}) \ (\text{inverse } (2 \wedge n)) = \text{global-depth-set}$
 (int n)
 ⟨proof⟩

lemma *p-adic-prod-metric-ball-UNIV*:

$\text{ball } x \ e = \text{UNIV}$ **if** $e > 1$
for $x :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$
 ⟨proof⟩

lemma *p-adic-prod-metric-ball-at-0-normalize*:

$\text{ball } (0::'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}) \ e =$
 $\text{global-depth-set } \lfloor \log 2 \ (\text{inverse } e) \rfloor$
if $e > 0$ **and** $e \leq 1$
 ⟨proof⟩

lemma *p-adic-prod-metric-ball-translate*:

$\text{ball } x \ e = x + o \ \text{ball } 0 \ e$ **for** $x :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$
 ⟨proof⟩

lemma *p-adic-prod-left-translate-metric-continuous*:

$\text{continuous-on UNIV } ((+) \ x)$ **for** $x :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$
 ⟨proof⟩

lemma *p-adic-prod-right-translate-metric-continuous*:

$\text{continuous-on UNIV } (\lambda z. z + x)$ **for** $x :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$
 ⟨proof⟩

lemma *p-adic-prod-left-bdd-mult-bdd-metric-continuous*:

$\text{continuous-on UNIV } ((*)) \ x$

if *bdd*:
 $\forall p::'a \text{ prime. } x \neg\simeq_p 0 \longrightarrow$
 $(x^{\circ p}) \geq n$
for $x :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-bdd-metric-limseq-bdd-mult*:
 $(\lambda k. w * X k) \longrightarrow (w * x)$
if *bdd*:
 $\forall p::'a \text{ prime. } w \neg\simeq_p 0 \longrightarrow$
 $(w^{\circ p}) \geq n$
and $\text{seq: } X \longrightarrow x$
for $w x :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$ **and** $X :: \text{nat} \Rightarrow 'a \text{ p-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-nonpos-global-depth-set-open*:
 $\text{ball } x \ 1 \subseteq (\mathcal{P}_{\forall p}^n)$
if $n \leq 0$ **and** $x \in (\mathcal{P}_{\forall p}^n)$
for $x :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-global-depth-set-open*:
 $\text{open } ((\mathcal{P}_{\forall p}^n) :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod set})$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-ball-abs-eq*:
 $\text{abs-p-adic-prod } \{ Y. \text{fls-pseq-bdd-global-dist } X \ Y < e \} =$
 $\text{ball } (\text{abs-p-adic-prod } X) \ e$
for $x :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

Completeness

abbreviation *p-adic-prod-globally-cauchy* $\equiv \text{global-p-depth-p-adic-prod.globally-cauchy}$
abbreviation
p-adic-prod-global-cauchy-condition $\equiv \text{global-p-depth-p-adic-prod.global-cauchy-condition}$

lemma *p-adic-prod-global-cauchy-condition-abs-eq*:
 $\text{p-adic-prod-global-cauchy-condition } (\lambda n. \text{abs-p-adic-prod } (X \ n)) =$
 $\text{fls-pseq-global-cauchy-condition } X$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-globally-cauchy-abs-eq*:
 $\text{p-adic-prod-globally-cauchy } (\lambda n. \text{abs-p-adic-prod } (X \ n)) = \text{fls-pseq-globally-cauchy}$
 X
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-globally-cauchy-vs-bdd-metric-Cauchy*:
 $\text{p-adic-prod-globally-cauchy } X = \text{Cauchy } X$
for $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$

$\langle \text{proof} \rangle$

abbreviation

$p\text{-adic-prod-cauchy-lim } X \equiv$
 $\text{abs-}p\text{-adic-prod } (\lambda p.$
 $\text{fls-condition-lim } p$
 $(\lambda n. \text{reduced-rep-}p\text{-adic-prod } (X \ n) \ p)$
 $(p\text{-adic-prod-global-cauchy-condition } X)$
 $)$

lemma $p\text{-adic-prod-bdd-metric-complete}'$:

$X \longrightarrow p\text{-adic-prod-cauchy-lim } X$ **if** $\text{Cauchy } X$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-bdd-metric-complete}$:

$\text{complete } (UNIV :: 'a::\text{nontrivial-euclidean-ring-cancel } p\text{-adic-prod set})$
 $\langle \text{proof} \rangle$

5.7 Hiding implementation details

lifting-update $p\text{-adic-prod.lifting}$

lifting-forget $p\text{-adic-prod.lifting}$

5.8 The subring of adelic integers

5.8.1 Type definition as a subtype of adeles

typedef (**overloaded**) $'a \text{ adelic-int} =$

$\mathcal{O}_{\forall p} :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod set}$
 $\langle \text{proof} \rangle$

lemma $\text{Abs-adelic-int-inverse}'$:

$\text{Rep-adelic-int } (\text{Abs-adelic-int } x) = x$ **if** $x \in \mathcal{O}_{\forall p}$
 $\langle \text{proof} \rangle$

lemma $\text{Abs-adelic-int-cases}$ [$\text{case-names Abs-adelic-int}$, $\text{cases type: adelic-int}$]:

P **if**
 $\bigwedge x. z = \text{Abs-adelic-int } x \implies$
 $x \in \mathcal{O}_{\forall p} \implies P$

$\langle \text{proof} \rangle$

lemmas $\text{two-Abs-adelic-int-cases} = \text{Abs-adelic-int-cases}[\text{case-product Abs-adelic-int-cases}]$

and $\text{three-Abs-adelic-int-cases} =$

$\text{Abs-adelic-int-cases}[\text{case-product Abs-adelic-int-cases}[\text{case-product Abs-adelic-int-cases}]]$

lemma $\text{adelic-int-seq-cases}$ [$\text{case-names Abs-adelic-int}$]:

fixes $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-factorial-idom adelic-int}$

obtains $F :: \text{nat} \Rightarrow 'a \text{ p-adic-prod}$

where $X = (\lambda n. \text{Abs-adelic-int } (F \ n))$

and $\text{range } F \subseteq \mathcal{O}_{\forall p}$

$\langle proof \rangle$

5.8.2 Algebraic instantiaions

instantiation *adelic-int* :: (*nontrivial-factorial-idom*) *comm-ring-1*
begin

definition $0 = \text{Abs-adelic-int } 0$

lemma *zero-adelic-int-rep-eq*: $\text{Rep-adelic-int } 0 = 0$
 $\langle proof \rangle$

definition $1 \equiv \text{Abs-adelic-int } 1$

lemma *one-adelic-int-rep-eq*: $\text{Rep-adelic-int } 1 = 1$
 $\langle proof \rangle$

definition $x + y = \text{Abs-adelic-int } (\text{Rep-adelic-int } x + \text{Rep-adelic-int } y)$

lemma *plus-adelic-int-rep-eq*: $\text{Rep-adelic-int } (a + b) = \text{Rep-adelic-int } a + \text{Rep-adelic-int } b$
 $\langle proof \rangle$

lemma *plus-adelic-int-abs-eq*:
 $x \in \mathcal{O}_{\forall p} \implies$
 $y \in \mathcal{O}_{\forall p} \implies$
 $\text{Abs-adelic-int } x + \text{Abs-adelic-int } y = \text{Abs-adelic-int } (x + y)$
 $\langle proof \rangle$

definition $-x = \text{Abs-adelic-int } (- \text{Rep-adelic-int } x)$

lemma *uminus-adelic-int-rep-eq*: $\text{Rep-adelic-int } (- x) = - \text{Rep-adelic-int } x$
 $\langle proof \rangle$

lemma *uminus-adelic-int-abs-eq*:
 $x \in \mathcal{O}_{\forall p} \implies - \text{Abs-adelic-int } x = \text{Abs-adelic-int } (- x)$
 $\langle proof \rangle$

definition *minus-adelic-int* ::
 $'a \text{ adelic-int} \Rightarrow 'a \text{ adelic-int} \Rightarrow 'a \text{ adelic-int}$
where $\text{minus-adelic-int} \equiv (\lambda x y. x + - y)$

lemma *minus-adelic-int-rep-eq*: $\text{Rep-adelic-int } (x - y) = \text{Rep-adelic-int } x - \text{Rep-adelic-int } y$
 $\langle proof \rangle$

lemma *minus-adelic-int-abs-eq*:
 $x \in \mathcal{O}_{\forall p} \implies$
 $y \in \mathcal{O}_{\forall p} \implies$

$Abs\text{-}adelic\text{-}int\ x - Abs\text{-}adelic\text{-}int\ y = Abs\text{-}adelic\text{-}int\ (x - y)$
 $\langle proof \rangle$

definition $x * y = Abs\text{-}adelic\text{-}int\ (Rep\text{-}adelic\text{-}int\ x * Rep\text{-}adelic\text{-}int\ y)$

lemma $times\text{-}adelic\text{-}int\text{-}rep\text{-}eq$: $Rep\text{-}adelic\text{-}int\ (x * y) = Rep\text{-}adelic\text{-}int\ x * Rep\text{-}adelic\text{-}int\ y$
 $\langle proof \rangle$

lemma $times\text{-}adelic\text{-}int\text{-}abs\text{-}eq$:
 $x \in \mathcal{O}_{\forall p} \implies$
 $y \in \mathcal{O}_{\forall p} \implies$
 $Abs\text{-}adelic\text{-}int\ x * Abs\text{-}adelic\text{-}int\ y = Abs\text{-}adelic\text{-}int\ (x * y)$
 $\langle proof \rangle$

instance
 $\langle proof \rangle$

end

lemma $pow\text{-}adelic\text{-}int\text{-}rep\text{-}eq$: $Rep\text{-}adelic\text{-}int\ (x \wedge n) = (Rep\text{-}adelic\text{-}int\ x) \wedge n$
 $\langle proof \rangle$

lemma $pow\text{-}adelic\text{-}int\text{-}abs\text{-}eq$:
 $(Abs\text{-}adelic\text{-}int\ x) \wedge n = Abs\text{-}adelic\text{-}int\ (x \wedge n)$ **if** $x \in \mathcal{O}_{\forall p}$
 $\langle proof \rangle$

5.8.3 Equivalence and depth relative to a prime

overloading

$p\text{-}equal\text{-}adelic\text{-}int \equiv$
 $p\text{-}equal :: 'a :: nontrivial\text{-}factorial\text{-}idom\ prime \Rightarrow$
 $'a\ adelic\text{-}int \Rightarrow 'a\ adelic\text{-}int \Rightarrow bool$
 $p\text{-}restrict\text{-}adelic\text{-}int \equiv$
 $p\text{-}restrict :: 'a\ adelic\text{-}int \Rightarrow$
 $('a\ prime \Rightarrow bool) \Rightarrow 'a\ adelic\text{-}int$
 $p\text{-}depth\text{-}adelic\text{-}int \equiv$
 $p\text{-}depth :: 'a :: nontrivial\text{-}factorial\text{-}idom\ prime \Rightarrow 'a\ adelic\text{-}int \Rightarrow int$
 $global\text{-}unfrmzr\text{-}pows\text{-}adelic\text{-}int \equiv$
 $global\text{-}unfrmzr\text{-}pows :: ('a\ prime \Rightarrow nat) \Rightarrow 'a\ adelic\text{-}int$

begin

definition $p\text{-}equal\text{-}adelic\text{-}int ::$
 $'a :: nontrivial\text{-}factorial\text{-}idom\ prime \Rightarrow$
 $'a\ adelic\text{-}int \Rightarrow 'a\ adelic\text{-}int \Rightarrow bool$
where $p\text{-}equal\text{-}adelic\text{-}int\ p\ x\ y \equiv (Rep\text{-}adelic\text{-}int\ x) \simeq_p (Rep\text{-}adelic\text{-}int\ y)$

definition $p\text{-}restrict\text{-}adelic\text{-}int ::$
 $'a :: nontrivial\text{-}factorial\text{-}idom\ adelic\text{-}int \Rightarrow$

$(\text{'a prime} \Rightarrow \text{bool}) \Rightarrow \text{'a adelic-int}$
where
 $p\text{-restrict-adelic-int } x \ P \equiv \text{Abs-adelic-int } ((\text{Rep-adelic-int } x) \text{ prestrict } P)$

definition $p\text{-depth-adelic-int} ::$
 $\text{'a} :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{'a adelic-int} \Rightarrow \text{int}$
where $p\text{-depth-adelic-int } p \ x \equiv ((\text{Rep-adelic-int } x)^{\circ p})$

definition $\text{global-unfrmzr-pows-adelic-int} ::$
 $(\text{'a} :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{nat}) \Rightarrow \text{'a adelic-int}$
where
 $\text{global-unfrmzr-pows-adelic-int } f \equiv \text{Abs-adelic-int } (\text{global-unfrmzr-pows } (\text{int} \circ f))$

end

context
fixes $p :: \text{'a} :: \text{nontrivial-factorial-idom prime}$
begin

lemma $p\text{-equal-adelic-int-abs-eq}$:
 $(\text{Abs-adelic-int } x) \simeq_p (\text{Abs-adelic-int } y) \longleftrightarrow$
 $x \simeq_p y$
if $x \in \mathcal{O}_{\forall p}$
and $y \in \mathcal{O}_{\forall p}$
for $x \ y :: \text{'a } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemma $p\text{-equal-adelic-int-abs-eq0}$:
 $(\text{Abs-adelic-int } x) \simeq_p 0 \longleftrightarrow x \simeq_p 0$
if $x \in \mathcal{O}_{\forall p}$ **for** $x :: \text{'a } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemma $p\text{-equal-adelic-int-rep-eq0}$:
 $(\text{Rep-adelic-int } x) \simeq_p 0 \longleftrightarrow x \simeq_p 0$
for $x :: \text{'a adelic-int}$
 $\langle \text{proof} \rangle$

lemma $p\text{-equal-adelic-int-abs-eq1}$:
 $(\text{Abs-adelic-int } x) \simeq_p 1 \longleftrightarrow x \simeq_p 1$
if $x \in \mathcal{O}_{\forall p}$ **for** $x :: \text{'a } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemma $p\text{-depth-adelic-int-abs-eq}$:
 $(\text{Abs-adelic-int } x)^{\circ p} = (x^{\circ p})$
if $x \in \mathcal{O}_{\forall p}$ **for** $x :: \text{'a } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemma adelic-int-depth : $(x^{\circ p}) \geq 0$ **for** $x :: \text{'a adelic-int}$

```

    <proof>

end

lemma p-restrict-adelic-int-rep-eq:
  Rep-adelic-int (x prestrict P) = (Rep-adelic-int x) prestrict P
  for x :: 'a::nontrivial-factorial-idom adelic-int
  and P :: 'a prime  $\Rightarrow$  bool
  <proof>

lemma p-restrict-adelic-int-abs-eq:
  (Abs-adelic-int x) prestrict P = Abs-adelic-int (x prestrict P)
  if  $x \in \mathcal{O}_{\forall p}$ 
  for x :: 'a::nontrivial-factorial-idom p-adic-prod and P :: 'a prime  $\Rightarrow$  bool
  <proof>

lemma global-unfrmzr-pows-adelic-int-rep-eq:
  Rep-adelic-int (p f :: 'a adelic-int) = p (int  $\circ$  f)
  for f :: 'a::nontrivial-factorial-idom prime  $\Rightarrow$  nat
  <proof>

global-interpretation p-equality-adelic-int:
  p-equality-no-zero-divisors-1
  p-equal :: 'a::nontrivial-factorial-idom prime  $\Rightarrow$ 
    'a adelic-int  $\Rightarrow$  'a adelic-int  $\Rightarrow$  bool
  <proof>

overloading
  globally-p-equal-adelic-int  $\equiv$ 
  globally-p-equal ::
    'a::nontrivial-factorial-idom adelic-int  $\Rightarrow$  'a adelic-int  $\Rightarrow$  bool
begin

definition globally-p-equal-adelic-int ::
  'a::nontrivial-factorial-idom adelic-int  $\Rightarrow$  'a adelic-int  $\Rightarrow$  bool
  where globally-p-equal-adelic-int-def[simp]:
    globally-p-equal-adelic-int = p-equality-adelic-int.globally-p-equal

end

global-interpretation global-p-depth-adelic-int:
  global-p-equal-depth-no-zero-divisors-w-inj-index-consts
  p-equal :: 'a::nontrivial-factorial-idom prime  $\Rightarrow$ 
    'a adelic-int  $\Rightarrow$  'a adelic-int  $\Rightarrow$  bool
  p-restrict ::
    'a adelic-int  $\Rightarrow$  ('a prime  $\Rightarrow$  bool)  $\Rightarrow$  'a adelic-int
  p-depth :: 'a prime  $\Rightarrow$  'a adelic-int  $\Rightarrow$  int
   $\lambda f$ . Abs-adelic-int (p-adic-prod-consts f)
  Rep-prime

```

$\langle proof \rangle$

lemma *p-depth-adelic-int-diff-abs-eq*:

$((Abs\text{-}adelic\text{-}int\ x - Abs\text{-}adelic\text{-}int\ y)^{\circ p}) = ((x - y)^{\circ p})$

if $x \in \mathcal{O}_{\mathbb{V}_p}$ **and** $y \in \mathcal{O}_{\mathbb{V}_p}$

for $p :: 'a::nontrivial\text{-}factorial\text{-}idom\ prime$ **and** $x\ y :: 'a\ p\text{-}adic\text{-}prod$

$\langle proof \rangle$

lemma *p-depth-adelic-int-diff-abs-eq'*:

$((Abs\text{-}adelic\text{-}int\ (abs\text{-}p\text{-}adic\text{-}prod\ X) - Abs\text{-}adelic\text{-}int\ (abs\text{-}p\text{-}adic\text{-}prod\ Y))^{\circ p} = ((X - Y)^{\circ p})$

if $X: X \in \mathcal{O}_{\mathbb{V}_p}$ **and** $Y: Y \in \mathcal{O}_{\mathbb{V}_p}$

for $p :: 'a::nontrivial\text{-}factorial\text{-}idom\ prime$ **and** $X\ Y :: 'a\ fls\text{-}pseq$

$\langle proof \rangle$

context

fixes $p :: 'a::nontrivial\text{-}euclidean\text{-}ring\text{-}cancel\ prime$

begin

lemma *adelic-int-diff-depth-gt-equiv-trans*:

$((a - c)^{\circ p}) > n$

if $a \simeq_p b$ **and** $b \neg\simeq_p c$ **and** $((b - c)^{\circ p}) > n$

for $a\ b\ c :: 'a\ adelic\text{-}int$

$\langle proof \rangle$

lemma *adelic-int-diff-depth-gt0-equiv-trans*:

$((a - c)^{\circ p}) > n$

if $n \geq 0$ **and** $a \simeq_p b$ **and** $((b - c)^{\circ p}) > n$

for $a\ b\ c :: 'a\ adelic\text{-}int$

$\langle proof \rangle$

lemma *adelic-int-diff-depth-gt-trans*:

$((a - c)^{\circ p}) > n$

if $a \neg\simeq_p b$ **and** $((a - b)^{\circ p}) > n$

and $b \neg\simeq_p c$ **and** $((b - c)^{\circ p}) > n$

and $a \neg\simeq_p c$

for $a\ b\ c :: 'a\ adelic\text{-}int$

$\langle proof \rangle$

lemma *adelic-int-diff-depth-gt0-trans*:

$((a - c)^{\circ p}) > n$

if $n \geq 0$

and $((a - b)^{\circ p}) > n$

and $((b - c)^{\circ p}) > n$

and $a \neg\simeq_p c$

for $a\ b\ c :: 'a\ adelic\text{-}int$

$\langle proof \rangle$

lemma *adelic-int-prestrict-zero-depth*:

```

1 ≤ ((x prestrict (λp::'a prime. (xop) = 0) - x)op)
if x prestrict (λp::'a prime. (xop) = 0) ¬p x
for x :: 'a adelic-int
⟨proof⟩

end

lemma adelic-int-normalize-depth:
  (x * p (λp::'a prime. -(xop))) ∈  $\mathcal{O}_{\forall p}$ 
for x :: 'a::nontrivial-factorial-idom p-adic-prod
⟨proof⟩

lemma global-unfrmzr-pows0-adelic-int:
  (p (0 :: 'a::nontrivial-factorial-idom prime ⇒ nat) :: 'a adelic-int) = 1
⟨proof⟩

context
  fixes p :: 'a::nontrivial-factorial-idom prime
begin

lemma global-unfrmzr-pows-adelic-int:
  (p f :: 'a adelic-int)op = int (f p) for f :: 'a prime ⇒ nat
⟨proof⟩

lemma global-unfrmzr-pows-adelic-int-pequiv-iff:
  ((p f :: 'a adelic-int) p (p g)) = (f p = g p)
for f g :: 'a prime ⇒ nat
⟨proof⟩

end

lemma global-unfrmzr-pows-prod-adelic-int:
  (p f :: 'a adelic-int) * (p g) = p (f + g)
for f g :: 'a::nontrivial-factorial-idom prime ⇒ nat
⟨proof⟩

context
  fixes p :: 'a::nontrivial-factorial-idom prime
begin

lemma global-unfrmzr-pows-adelic-int-nzero:
  (p f :: 'a adelic-int) ¬p 0 for f :: 'a prime ⇒ nat
⟨proof⟩

lemma prod-w-global-unfrmzr-pows-adelic-int:
  (x * p f)op = (xop) + int (f p)
if x ¬p 0 for f :: 'a prime ⇒ nat and x :: 'a adelic-int
⟨proof⟩

```

lemma *trivial-global-unfrmzr-pows-adelic-int*:
 $(\mathbf{p} \ f :: 'a \text{ adelic-int}) \simeq_p 1$ **if** $f \ p = 0$ **for** $f :: 'a \text{ prime} \Rightarrow \text{nat}$
 $\langle \text{proof} \rangle$

lemma *prod-w-trivial-global-unfrmzr-pows-adelic-int*:
 $x * \mathbf{p} \ f \simeq_p x$ **if** $f \ p = 0$
for $f :: 'a \text{ prime} \Rightarrow \text{nat}$ **and** $x :: 'a \text{ adelic-int}$
 $\langle \text{proof} \rangle$

end

lemma *pow-global-unfrmzr-pows-adelic-int*:
 $(\mathbf{p} \ f :: 'a \text{ adelic-int}) \wedge^n = (\mathbf{p} \ ((\lambda p. n) * f))$
for $f :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{nat}$
 $\langle \text{proof} \rangle$

abbreviation *adelic-int-consts* $\equiv$ *global-p-depth-adelic-int.consts*

abbreviation *adelic-int-const* $\equiv$ *global-p-depth-adelic-int.const*

lemma *adelic-int-p-depth-consts*:
 $(\text{adelic-int-consts } f)^{\circ p} = ((f \ p)^{\circ p})$
for $p :: 'a :: \text{nontrivial-factorial-idom prime}$ **and** $f :: 'a \text{ prime} \Rightarrow 'a$
 $\langle \text{proof} \rangle$

lemmas *adelic-int-p-depth-const* $=$ *adelic-int-p-depth-consts*[*of* - $\lambda p. -$]

lemma *adelic-int-global-unfrmzr*:
 $(\mathbf{p} \ (1 :: 'a :: \text{nontrivial-factorial-idom prime} \Rightarrow \text{nat}) :: 'a \text{ adelic-int})$
 $=$ *adelic-int-consts Rep-prime*
 $\langle \text{proof} \rangle$

lemma *adelic-int-normalize-depthE*:
fixes $x :: 'a :: \text{nontrivial-factorial-idom adelic-int}$
obtains $x \cdot 0$
where $\forall p :: 'a \text{ prime. } x \cdot 0^{\circ p} = 0$
and $x = x \cdot 0 * \mathbf{p} \ (\lambda p :: 'a \text{ prime. nat } (x^{\circ p}))$
 $\langle \text{proof} \rangle$

5.8.4 Inverses

We can create inverses at depth-zero places.

instantiation *adelic-int* ::
 $(\text{nontrivial-factorial-unique-euclidean-bezout}) \{ \text{divide-trivial, inverse} \}$
begin

definition *inverse-adelic-int* ::
 $'a \text{ adelic-int} \Rightarrow 'a \text{ adelic-int}$
where
inverse-adelic-int $x =$

```

    Abs-adelic-int (
      (inverse (Rep-adelic-int x)) prestrict (λp::'a prime. xop = 0)
    )

```

definition *divide-adelic-int* ::
 'a adelic-int ⇒ 'a adelic-int ⇒ 'a adelic-int
where *divide-adelic-int-def[simp]*: *divide-adelic-int* x y = x * inverse y

instance

⟨proof⟩

end

lemma *inverse-adelic-int-abs-eq*:

```

inverse (Abs-adelic-int x) =
  Abs-adelic-int (inverse x prestrict (λp::'a prime. xop = 0))
if x ∈ O∇p
for x :: 'a::nontrivial-factorial-unique-euclidean-bezout p-adic-prod
⟨proof⟩

```

lemma *divide-adelic-int-abs-eq*:

```

Abs-adelic-int x / Abs-adelic-int y =
  Abs-adelic-int ((x / y) prestrict (λp::'a prime. yop = 0))
if x ∈ O∇p and y ∈ O∇p
for x y :: 'a::nontrivial-factorial-unique-euclidean-bezout p-adic-prod
⟨proof⟩

```

lemma *p-depth-adelic-int-inverse-diff-abs-eq*:

```

((inverse (Abs-adelic-int x) - inverse (Abs-adelic-int y))op) =
  ((inverse x - inverse y)op)
if x ∈ O∇p xop = 0
and y ∈ O∇p yop = 0
for p :: 'a::nontrivial-factorial-unique-euclidean-bezout prime and x y :: 'a p-adic-prod
⟨proof⟩

```

context

fixes x :: 'a::nontrivial-factorial-unique-euclidean-bezout adelic-int

begin

lemma *inverse-adelic-int-rep-eq*:

```

Rep-adelic-int (inverse x) =
  (inverse (Rep-adelic-int x)) prestrict (λp::'a prime. xop = 0)
⟨proof⟩

```

lemma *adelic-int-inverse-equiv0-iff*:

```

inverse x ≃p 0 ⟷
  (x ≃p 0 ∨ (xop) ≠ 0)
for p :: 'a prime
⟨proof⟩

```

lemma *adelic-int-inverse-equiv0-iff'*:

$inverse\ x \simeq_p 0 \longleftrightarrow$
 $(x \simeq_p 0 \vee (x^{op} \geq 1))$
for $p :: 'a\ prime$
 $\langle proof \rangle$

end

lemma *adelic-int-inverse-eq0-iff*:

$inverse\ x = 0 \longleftrightarrow$
 $(\forall p :: 'a\ prime. x^{op} = 0 \longrightarrow x \simeq_p 0)$
for $x :: 'a :: nontrivial-factorial-unique-euclidean-bezout\ adelic-int$
 $\langle proof \rangle$

lemma *divide-adelic-int-rep-eq*:

$Rep-adelic-int\ (x / y) =$
 $(Rep-adelic-int\ x / Rep-adelic-int\ y)\ prestrict\ (\lambda p :: 'a\ prime. y^{op} = 0)$
for $x\ y :: 'a :: nontrivial-factorial-unique-euclidean-bezout\ adelic-int$
 $\langle proof \rangle$

lemma *adelic-int-divide-by-pequiv0*:

$x / y \simeq_p 0$ **if** $y \simeq_p 0$
for $p :: 'a :: nontrivial-factorial-unique-euclidean-bezout\ prime$ **and** $x\ y :: 'a\ adelic-int$
 $\langle proof \rangle$

lemma *adelic-int-divide-by-pos-depth*:

$x / y \simeq_p 0$ **if** $(y^{op}) > 0$
for $p :: 'a :: nontrivial-factorial-unique-euclidean-bezout\ prime$ **and** $x\ y :: 'a\ adelic-int$
 $\langle proof \rangle$

lemma *adelic-int-inverse0[simp]*:

$inverse\ (0 :: 'a :: nontrivial-factorial-unique-euclidean-bezout\ adelic-int) = 0$
 $\langle proof \rangle$

lemma *adelic-int-inverse1[simp]*:

$inverse\ (1 :: 'a :: nontrivial-factorial-unique-euclidean-bezout\ adelic-int) = 1$
 $\langle proof \rangle$

lemma *adelic-int-inverse-mult*:

$inverse\ (x * y) = inverse\ x * inverse\ y$
for $x\ y :: 'a :: nontrivial-factorial-unique-euclidean-bezout\ adelic-int$
 $\langle proof \rangle$

lemma *adelic-int-inverse-pcong*:

$(inverse\ x) \simeq_p (inverse\ y)$ **if** $x \simeq_p y$
for $p :: 'a :: nontrivial-factorial-unique-euclidean-bezout\ prime$ **and** $x\ y :: 'a\ adelic-int$
 $\langle proof \rangle$

```

context
  fixes  $x :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout adelic-int}$ 
begin

lemma adelic-int-inverse-uminus:  $\text{inverse } (-x) = - \text{inverse } x$ 
  <proof>

lemma adelic-int-inverse-inverse:
   $\text{inverse } (\text{inverse } x) = x \text{ prestrict } (\lambda p :: 'a \text{ prime. } x^{\circ p} = 0)$ 
  <proof>

lemma adelic-int-inverse-pow:  $\text{inverse } (x \wedge n) = (\text{inverse } x) \wedge n$ 
  <proof>

lemma adelic-int-divide-self':
   $x / x \simeq_p 1$  if  $x \not\simeq_p 0$  and  $x^{\circ p} = 0$ 
  for  $p :: 'a \text{ prime}$ 
  <proof>

lemma adelic-int-global-divide-self:
   $x / x = 1$ 
  if  $\forall p :: 'a \text{ prime. } x \not\simeq_p 0$ 
  and  $\forall p :: 'a \text{ prime. } x^{\circ p} = 0$ 
  <proof>

lemma adelic-int-divide-self:
   $x / x =$ 
   $1 \text{ prestrict } (\lambda p :: 'a \text{ prime. } x \not\simeq_p 0 \wedge x^{\circ p} = 0)$ 
  <proof>

lemma adelic-int-p-restrict-inverse:
   $(\text{inverse } x) \text{ prestrict } P = \text{inverse } (x \text{ prestrict } P)$  for  $P :: 'a \text{ prime} \Rightarrow \text{bool}$ 
  <proof>

lemma adelic-int-inverse-depth:  $(\text{inverse } x)^{\circ p} = 0$  for  $p :: 'a \text{ prime}$ 
  <proof>

end

lemma adelic-int-consts-divide-self:
   $(\text{adelic-int-consts } f) / (\text{adelic-int-consts } f) =$ 
   $1 \text{ prestrict } (\lambda p :: 'a \text{ prime. } f \ p \neq 0 \wedge (f \ p)^{\circ p} = 0)$ 
  for  $f :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow 'a$ 
  <proof>

lemma adelic-int-const-divide-self:
   $(\text{adelic-int-const } c) / (\text{adelic-int-const } c) =$ 
   $1 \text{ prestrict } (\lambda p :: 'a \text{ prime. } c^{\circ p} = 0)$ 
  if  $c \neq 0$  for  $c :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout}$ 

```

$\langle \text{proof} \rangle$

lemma *adelic-int-consts-divide-self'*:

$(\text{adelic-int-consts } f) / (\text{adelic-int-consts } f) \simeq_p 1$

if $f p \neq 0$ **and** $(f p)^{\circ p} = 0$

for $f :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow 'a$

$\langle \text{proof} \rangle$

lemma *adelic-int-divide-depth*:

$(x / y)^{\circ p} = (x^{\circ p})$ **if** $x / y \neg\simeq_p 0$

for $x y :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout adelic-int}$ **and** $p :: 'a$
 prime

$\langle \text{proof} \rangle$

lemma *global-unfrmzr-pows-adelic-int-inverse*:

$\text{inverse } (p f :: 'a \text{ adelic-int}) = 1 \text{ prestrict } (\lambda p::'a \text{ prime. } f p = 0)$

for $f :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow \text{nat}$

$\langle \text{proof} \rangle$

lemma *adelic-int-diff-cancel-lead-coeff*:

$((\text{inverse } x - \text{inverse } y)^{\circ p}) > 0$

if $x : x \neg\simeq_p 0$ $x^{\circ p} = 0$

and $y : y \neg\simeq_p 0$ $y^{\circ p} = 0$

and $xy: x \neg\simeq_p y$ $((x - y)^{\circ p}) > 0$

for $p :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout prime}$

and $x y :: 'a \text{ adelic-int}$

$\langle \text{proof} \rangle$

5.8.5 The ideal of primes

overloading

$p\text{-depth-set-adelic-int} \equiv$

$p\text{-depth-set} ::$

$'a::\text{nontrivial-factorial-idom prime} \Rightarrow \text{int} \Rightarrow 'a \text{ adelic-int set}$

$\text{global-depth-set-adelic-int} \equiv$

$\text{global-depth-set} :: \text{int} \Rightarrow 'a \text{ adelic-int set}$

begin

definition *p-depth-set-adelic-int* ::

$'a::\text{nontrivial-factorial-idom prime} \Rightarrow \text{int} \Rightarrow 'a \text{ adelic-int set}$

where $p\text{-depth-set-adelic-int-def}[simp]$:

$p\text{-depth-set-adelic-int} = \text{global-p-depth-adelic-int}.p\text{-depth-set}$

definition *global-depth-set-adelic-int* ::

$\text{int} \Rightarrow 'a::\text{nontrivial-factorial-idom adelic-int set}$

where $\text{global-depth-set-adelic-int-def}[simp]$:

$\text{global-depth-set-adelic-int} = \text{global-p-depth-adelic-int}.global\text{-depth-set}$

end

```

context
  fixes  $x :: 'a::\text{nontrivial-factorial-idom adelic-int}$ 
begin

lemma adelic-int-global-depth-setI:
   $x \in \mathcal{P}_{\forall p}^n$ 
  if  $\bigwedge p::'a \text{ prime. } x \dashv\sim_p 0 \implies (x^{\circ p}) \geq n$ 
   $\langle \text{proof} \rangle$ 

lemma adelic-int-global-depth-setI':
   $x \in \mathcal{P}_{\forall p}^n$ 
  if  $\bigwedge p::'a \text{ prime. } (x^{\circ p}) \geq n$ 
   $\langle \text{proof} \rangle$ 

lemma adelic-int-global-depth-set-p-restrictI:
   $p\text{-restrict } x \ P \in \mathcal{P}_{\forall p}^n$ 
  if  $\bigwedge p::'a \text{ prime. } P \ p \implies x \in \mathcal{P}_{@p}^n$ 
   $\langle \text{proof} \rangle$ 

end

context
  fixes  $p :: 'a::\text{nontrivial-factorial-idom prime}$ 
begin

lemma adelic-int-p-depth-setI:
   $x \in \mathcal{P}_{@p}^n$ 
  if  $x \dashv\sim_p 0 \implies (x^{\circ p}) \geq n$ 
  for  $x :: 'a \text{ adelic-int}$ 
   $\langle \text{proof} \rangle$ 

lemma nonpos-p-depth-set-adelic-int:
   $(\mathcal{P}_{@p}^n) = (\text{UNIV} :: 'a::\text{nontrivial-factorial-idom adelic-int set})$ 
  if  $n \leq 0$ 
   $\langle \text{proof} \rangle$ 

lemma p-depth-set-adelic-int-eq-projection:
   $((\mathcal{P}_{@p}^n) :: 'a::\text{nontrivial-factorial-idom adelic-int set}) =$ 
   $\text{Abs-adelic-int } ((\mathcal{P}_{@p}^n) \cap \mathcal{O}_{\forall p})$ 
   $\langle \text{proof} \rangle$ 

lemma lift-p-depth-set-adelic-int:
   $\text{Rep-adelic-int } ((\mathcal{P}_{@p}^n) :: 'a::\text{nontrivial-factorial-idom adelic-int set}) =$ 
   $(\mathcal{P}_{@p}^n) \cap \mathcal{O}_{\forall p}$ 
   $\langle \text{proof} \rangle$ 

end

```

lemma *nonpos-global-depth-set-adelic-int*:
 $(\mathcal{P}_{\forall p}^n) = (\text{UNIV} :: 'a::\text{nontrivial-factorial-idom adelic-int set})$
if $n \leq 0$
 $\langle \text{proof} \rangle$

lemma *nonneg-global-depth-set-adelic-int-eq-projection*:
 $((\mathcal{P}_{\forall p}^n) :: 'a::\text{nontrivial-factorial-idom adelic-int set}) =$
 $\text{Abs-adelic-int } \cdot \mathcal{P}_{\forall p}^n$
if $n: n \geq 0$
 $\langle \text{proof} \rangle$

lemma *lift-nonneg-global-depth-set-adelic-int*:
 $\text{Rep-adelic-int } \cdot$
 $((\mathcal{P}_{\forall p}^n) :: 'a::\text{nontrivial-factorial-idom adelic-int set})$
 $= \mathcal{P}_{\forall p}^n$
if $n: n \geq 0$
 $\langle \text{proof} \rangle$

5.8.6 Topology p-open neighbourhoods

General properties of p-open sets

abbreviation *adelic-int-p-open-nbhd* $\equiv \text{global-p-depth-adelic-int.p-open-nbhd}$

abbreviation *adelic-int-p-open-nbhds* $\equiv \text{global-p-depth-adelic-int.p-open-nbhds}$

abbreviation *adelic-int-local-p-open-nbhds* $\equiv \text{global-p-depth-adelic-int.local-p-open-nbhds}$

lemma *adelic-int-nonpos-p-open-nbhd*:
 $\text{adelic-int-p-open-nbhd } p \ n \ x = \text{UNIV}$ **if** $n \leq 0$
for $p :: 'a::\text{nontrivial-factorial-idom prime}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-p-open-nbhd-bound*:
 $\text{adelic-int-p-open-nbhd } p \ n \ x = \text{adelic-int-p-open-nbhd } p \ (\text{int } (\text{nat } n)) \ x$
 $\langle \text{proof} \rangle$

lemma *adelic-int-p-open-nbhd-abs-eq*:
 $\text{adelic-int-p-open-nbhd } p \ n \ (\text{Abs-adelic-int } x) =$
 $\text{Abs-adelic-int } \cdot (p\text{-adic-prod-p-open-nbhd } p \ n \ x \cap \mathcal{O}_{\forall p})$
if $x \in \mathcal{O}_{\forall p}$
for $p :: 'a::\text{nontrivial-factorial-idom prime}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-p-open-nbhd-rep-eq*:
 $\text{Rep-adelic-int } \cdot \text{adelic-int-p-open-nbhd } p \ n \ x =$
 $p\text{-adic-prod-p-open-nbhd } p \ n \ (\text{Rep-adelic-int } x) \cap \mathcal{O}_{\forall p}$
for $p :: 'a::\text{nontrivial-factorial-idom prime}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-local-depth-ring-lift-cover*:

```

fixes  $n :: \text{int}$ 
and  $p :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime}$ 
and  $\mathcal{A} :: 'a \text{ adelic-int set set}$ 
defines
  ( $\mathcal{A}' :: 'a \text{ p-adic-prod set set}$ )  $\equiv$ 
  ( $\lambda A.$ 
    ( $\lambda x. x \text{ prestrict } ((=) p)$ ) '  $\text{Rep-adelic-int } A +$ 
     $\text{range } (\lambda x. x \text{ prestrict } ((\neq) p))$ 
  ) '  $\mathcal{A}$ 
assumes  $\text{nonneg: } n \geq 0$ 
and  $\text{cover:}$ 
  ( $\lambda x. x \text{ prestrict } ((=) p)$ ) ' ( $\mathcal{P}_{\forall p}^n$ )  $\subseteq$ 
   $\bigcup \mathcal{A}$ 
shows
  ( $\lambda x. x \text{ prestrict } ((=) p)$ ) ' ( $\mathcal{P}_{\forall p}^n$ )  $\subseteq$ 
   $\bigcup \mathcal{A}'$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma  $\text{adelic-int-lift-local-p-open:}$ 
fixes  $p :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout prime}$ 
and  $A :: 'a \text{ adelic-int set}$ 
defines
   $A' \equiv$ 
  ( $\lambda x. x \text{ prestrict } ((=) p)$ ) '  $\text{Rep-adelic-int } A +$ 
   $\text{range } (\lambda x. x \text{ prestrict } ((\neq) p))$ 
assumes  $\text{adelic-int-p-open: generate-topology (adelic-int-local-p-open-nbhds } p) A$ 
shows  $\text{generate-topology (p-adic-prod-local-p-open-nbhds } p) A'$ 
 $\langle \text{proof} \rangle$ 

```

Sequences

abbreviation

$\text{adelic-int-p-limseq-condition} \equiv \text{global-p-depth-adelic-int.p-limseq-condition}$

abbreviation

$\text{adelic-int-p-cauchy-condition} \equiv \text{global-p-depth-adelic-int.p-cauchy-condition}$

abbreviation $\text{adelic-int-p-limseq} \equiv \text{global-p-depth-adelic-int.p-limseq}$

abbreviation $\text{adelic-int-p-cauchy} \equiv \text{global-p-depth-adelic-int.p-cauchy}$

abbreviation $\text{adelic-int-p-open-nbhds-limseq} \equiv \text{global-p-depth-adelic-int.p-open-nbhds-LIMSEQ}$

context

fixes $p :: 'a :: \text{nontrivial-factorial-idom prime}$

and $X :: \text{nat} \Rightarrow 'a \text{ p-adic-prod}$

and $Y :: \text{nat} \Rightarrow 'a \text{ adelic-int}$

defines $Y \equiv \lambda n. \text{Abs-adelic-int } (X n)$

assumes $\text{range-X: range } X \subseteq \mathcal{O}_{\forall p}$

begin

lemma $\text{adelic-int-p-limseq-condition-abs-eq:}$

$\text{adelic-int-p-limseq-condition } p Y (\text{Abs-adelic-int } x) =$
 $\text{p-adic-prod-p-limseq-condition } p X x$

if $x \in \mathcal{O}_{\forall p}$
 $\langle proof \rangle$

lemma *adelic-int-p-limseq-abs-eq*:
 $adelic-int-p-limseq\ p\ Y\ (Abs-adelic-int\ x) = p-adic-prod-p-limseq\ p\ X\ x$
 if $x \in \mathcal{O}_{\forall p}$
 $\langle proof \rangle$

lemma *adelic-int-p-cauchy-condition-abs-eq*:
 $adelic-int-p-cauchy-condition\ p\ Y = p-adic-prod-p-cauchy-condition\ p\ X$
 $\langle proof \rangle$

lemma *adelic-int-p-cauchy-abs-eq*: $adelic-int-p-cauchy\ p\ Y = p-adic-prod-p-cauchy\ p\ X$
 $\langle proof \rangle$

end

lemma *adelic-int-p-open-nbhds-limseq-abs-eq*:
 $adelic-int-p-open-nbhds-limseq\ (\lambda n. Abs-adelic-int\ (X\ n))\ (Abs-adelic-int\ x) =$
 $p-adic-prod-p-open-nbhds-limseq\ X\ x$
 if $range\ X \subseteq \mathcal{O}_{\forall p}$ and $x \in \mathcal{O}_{\forall p}$
 $\langle proof \rangle$

5.8.7 Topology via global depth

The metric

instantiation *adelic-int* :: (nontrivial-factorial-idom) metric-space
begin

definition *dist* = *global-p-depth-adelic-int.bdd-global-dist*
declare *dist-adelic-int-def* [simp]

definition *uniformity* = *global-p-depth-adelic-int.bdd-global-uniformity*
declare *uniformity-adelic-int-def* [simp]

definition *open-adelic-int* :: 'a *adelic-int set* $\Rightarrow$ bool
where
 $open-adelic-int\ U =$
 $(\forall x \in U.$
 $eventually\ (\lambda(x', y). x' = x \longrightarrow y \in U)\ uniformity$
 $)$

instance
 $\langle proof \rangle$

end

abbreviation *adelic-int-global-depth* $\equiv$ *global-p-depth-adelic-int.bdd-global-depth*

lemma *adelic-int-bdd-global-depth-abs-eq*:
 $\text{adelic-int-global-depth } (\text{Abs-adelic-int } x) = \text{p-adic-prod-global-depth } x$
if $x \in \mathcal{O}_{\mathbb{V}_p}$ **for** $x :: 'a::\text{nontrivial-factorial-idom p-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *dist-adelic-int-abs-eq*:
 $\text{dist } (\text{Abs-adelic-int } x) (\text{Abs-adelic-int } y) = \text{dist } x y$
if $x \in \mathcal{O}_{\mathbb{V}_p}$ **and** $y \in \mathcal{O}_{\mathbb{V}_p}$
for $x y :: 'a::\text{nontrivial-factorial-idom p-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-limseq-abs-eq*:
 $(\lambda n. \text{Abs-adelic-int } (X n)) \longrightarrow \text{Abs-adelic-int } x$
 $\longleftrightarrow X \longrightarrow x$
if $\text{range } X \subseteq \mathcal{O}_{\mathbb{V}_p}$ **and** $x \in \mathcal{O}_{\mathbb{V}_p}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-bdd-metric-LIMSEQ*:
 $X \longrightarrow x = \text{global-p-depth-adelic-int.metric-space-by-bdd-depth.LIMSEQ } X x$
for $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-factorial-idom adelic-int}$ **and** $x :: 'a \text{ adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-global-depth-set-lim-closed*:
 $x \in \mathcal{P}_{\mathbb{V}_p}^n$
if $\text{lim} : X \longrightarrow x$
and $\text{depth} : \forall_F k \text{ in sequentially. } X k \in \mathcal{P}_{\mathbb{V}_p}^n$
for $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-factorial-unique-euclidean-bezout adelic-int}$
and $x :: 'a \text{ adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-metric-ball-multicentre*:
 $\text{ball } y e = \text{ball } x e$ **if** $y \in \text{ball } x e$ **for** $x y :: 'a::\text{nontrivial-factorial-idom adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-metric-ball-at-0*:
 $\text{ball } (0 :: 'a::\text{nontrivial-factorial-idom adelic-int}) (\text{inverse } (2 \wedge n)) = \text{global-depth-set } (int n)$
 $\langle \text{proof} \rangle$

lemma *adelic-int-metric-ball-at-0-normalize*:
 $\text{ball } (0 :: 'a::\text{nontrivial-factorial-idom adelic-int}) e =$
 $\text{global-depth-set } \lfloor \log 2 (\text{inverse } e) \rfloor$
if $e > 0$ **and** $e \leq 1$
 $\langle \text{proof} \rangle$

lemma *adelic-int-metric-ball-translate*:
 $\text{ball } x e = x +_o \text{ball } 0 e$ **for** $x :: 'a::\text{nontrivial-factorial-idom adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-metric-ball-UNIV*:

ball x $e = \text{UNIV}$ **if** $e \geq 1$ **for** $x :: 'a::\text{nontrivial-factorial-idom adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-left-translate-metric-continuous*:

continuous-on UNIV $((+) x)$ **for** $x :: 'a::\text{nontrivial-factorial-idom adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-right-translate-metric-continuous*:

continuous-on UNIV $(\lambda z. z + x)$ **for** $x :: 'a::\text{nontrivial-factorial-idom adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-left-mult-bdd-metric-continuous*:

continuous-on UNIV $((*) x)$
for $x :: 'a::\text{nontrivial-factorial-idom adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-bdd-metric-limseq-mult*:

$(\lambda k. w * X k) \longrightarrow (w * x)$ **if** $X \longrightarrow x$
for $w x :: 'a::\text{nontrivial-factorial-idom adelic-int}$ **and** $X :: \text{nat} \Rightarrow 'a \text{ adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-global-depth-set-open*:

open $((\mathcal{P}_{\forall p}^n) :: 'a::\text{nontrivial-factorial-idom adelic-int set})$
 $\langle \text{proof} \rangle$

lemma *adelic-int-ball-abs-eq*:

Abs-adelic-int ' *ball* x $e = \text{ball}$ (*Abs-adelic-int* x) e
if $x \in \mathcal{O}_{\forall p}$ **and** $e \leq 1$
for $x :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod}$
 $\langle \text{proof} \rangle$

lemma *open-adelic-int-abs-eq*:

open $U = \text{open}$ (*Abs-adelic-int* ' U) **if** $U \subseteq (\mathcal{O}_{\forall p})$
for $U :: 'a::\text{nontrivial-factorial-idom } p\text{-adic-prod set}$
 $\langle \text{proof} \rangle$

Completeness

abbreviation *adelic-int-globally-cauchy* $\equiv \text{global-}p\text{-depth-adelic-int.globally-cauchy}$

abbreviation

adelic-int-global-cauchy-condition $\equiv \text{global-}p\text{-depth-adelic-int.global-cauchy-condition}$

lemma *adelic-int-global-cauchy-condition-abs-eq*:

adelic-int-global-cauchy-condition $(\lambda n. \text{Abs-adelic-int } (X n)) =$
p-adic-prod-global-cauchy-condition X
if $\text{range } X \subseteq \mathcal{O}_{\forall p}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-globally-cauchy-abs-eq*:
adelic-int-globally-cauchy ($\lambda n. \text{Abs-adelic-int } (X \ n)$) = *p-adic-prod-globally-cauchy*
 X
if *range* $X \subseteq \mathcal{O}_{\forall p}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-globally-cauchy-vs-bdd-metric-Cauchy*:
adelic-int-globally-cauchy $X = \text{Cauchy } X$
for $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-factorial-idom adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-Cauchy-abs-eq*:
 $\text{Cauchy } (\lambda n. \text{Abs-adelic-int } (X \ n)) = \text{Cauchy } X$
if *range* $X \subseteq \mathcal{O}_{\forall p}$
for $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-factorial-idom p-adic-prod}$
 $\langle \text{proof} \rangle$

abbreviation
adelic-int-cauchy-lim $X \equiv$
 $\text{Abs-adelic-int } (p\text{-adic-prod-cauchy-lim } (\lambda n. \text{Rep-adelic-int } (X \ n)))$

lemma *adelic-int-bdd-metric-complete'*:
 $X \longrightarrow \text{adelic-int-cauchy-lim } X$ **if** *Cauchy* X
for $X :: \text{nat} \Rightarrow 'a::\text{nontrivial-factorial-unique-euclidean-bezout adelic-int}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-bdd-metric-complete*:
 $\text{complete } (\text{UNIV} :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout adelic-int set})$
 $\langle \text{proof} \rangle$

end

theory *Fin-Field-Product*

imports
Distinguished-Quotients
pAdic-Product

begin

6 Finite fields of prime order as quotients of the ring of integers of p-adic fields

6.1 The type

6.1.1 The prime ideal as the distinguished ideal of the ring of adelic integers

instantiation *adelic-int* ::
 (nontrivial-factorial-idom) comm-ring-1-with-distinguished-ideal
begin

definition *distinguished-subset-adelic-int* :: 'a adelic-int set
where *distinguished-subset-adelic-int-def*[simp]:
distinguished-subset-adelic-int $\equiv \mathcal{P}_{\mathbb{V}_p}$

instance
 <proof>

end

lemma *distinguished-subset-adelic-int-inverse*:
inverse $x = 0$ **if** $x \in \mathcal{P}_{\mathbb{V}_p}$
for $x :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout}$ *adelic-int*
 <proof>

6.1.2 The finite field product type as a quotient

Create type wrapper *coset-by-dist-sbgrp* over *adelic-int*.

typedef (overloaded) 'a *adelic-int-quot* =
 UNIV :: 'a::nontrivial-factorial-idom *adelic-int* *coset-by-dist-sbgrp* set
 <proof>

abbreviation

adelic-int-quot $\equiv \lambda x. \text{Abs-adelic-int-quot } (\text{distinguished-quotient-coset } x)$

abbreviation *p-adic-prod-int-quot* $\equiv \lambda x. \text{adelic-int-quot } (\text{Abs-adelic-int } x)$

lemma *adelic-int-quot-cases* [case-names *adelic-int-quot*]:
obtains y **where** $x = \text{adelic-int-quot } y$
 <proof>

lemma *Abs-adelic-int-quot-cases* [case-names *adelic-int-quot-Abs-adelic-int*]:
obtains z **where** $x = \text{p-adic-prod-int-quot } z$ $z \in \mathcal{O}_{\mathbb{V}_p}$
 <proof>

lemmas *two-adelic-int-quot-cases* = *adelic-int-quot-cases*[case-product *adelic-int-quot-cases*]
and *three-adelic-int-quot-cases* =
adelic-int-quot-cases[case-product *adelic-int-quot-cases*[case-product *adelic-int-quot-cases*]]

lemma *abs-adelic-int-eq-iff*:
 $\text{adelic-int-quot } x = \text{adelic-int-quot } y \longleftrightarrow$
 $y - x \in \mathcal{P}_{\mathbb{V}_p}$
 $\langle \text{proof} \rangle$

lemma *p-adic-prod-int-quot-eq-iff*:
 $p\text{-adic-prod-int-quot } x = p\text{-adic-prod-int-quot } y \longleftrightarrow$
 $y - x \in \mathcal{P}_{\mathbb{V}_p}$
if $x \in \mathcal{O}_{\mathbb{V}_p}$ **and** $y \in \mathcal{O}_{\mathbb{V}_p}$
 $\langle \text{proof} \rangle$

6.1.3 Algebraic instantiations

instantiation *adelic-int-quot* :: (*nontrivial-factorial-idom*) *comm-ring-1*
begin

definition $0 = \text{Abs-adelic-int-quot } 0$

lemma *zero-adelic-int-quot-rep-eq*: $\text{Rep-adelic-int-quot } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *zero-adelic-int-quot*: $0 = \text{adelic-int-quot } 0$
 $\langle \text{proof} \rangle$

definition $1 \equiv \text{Abs-adelic-int-quot } 1$

lemma *one-adelic-int-quot-rep-eq*: $\text{Rep-adelic-int-quot } 1 = 1$
 $\langle \text{proof} \rangle$

lemma *one-adelic-int-quot*: $1 = \text{adelic-int-quot } 1$
 $\langle \text{proof} \rangle$

definition $x + y = \text{Abs-adelic-int-quot } (\text{Rep-adelic-int-quot } x + \text{Rep-adelic-int-quot } y)$

lemma *plus-adelic-int-quot-rep-eq*:
 $\text{Rep-adelic-int-quot } (a + b) = \text{Rep-adelic-int-quot } a + \text{Rep-adelic-int-quot } b$
 $\langle \text{proof} \rangle$

lemma *plus-adelic-int-quot-abs-eq*:
 $\text{Abs-adelic-int-quot } x + \text{Abs-adelic-int-quot } y = \text{Abs-adelic-int-quot } (x + y)$
 $\langle \text{proof} \rangle$

lemma *plus-adelic-int-quot*: $\text{adelic-int-quot } x + \text{adelic-int-quot } y = \text{adelic-int-quot } (x + y)$
 $\langle \text{proof} \rangle$

definition $-x = \text{Abs-adelic-int-quot } (- \text{Rep-adelic-int-quot } x)$

lemma *uminus-adelic-int-quot-rep-eq*: $\text{Rep-adelic-int-quot } (-x) = - \text{Rep-adelic-int-quot } x$

<proof>

lemma *uminus-adelic-int-quot-abs-eq*: $- \text{Abs-adelic-int-quot } x = \text{Abs-adelic-int-quot } (-x)$

<proof>

lemma *uminus-adelic-int-quot*: $- \text{adelic-int-quot } x = \text{adelic-int-quot } (-x)$

<proof>

definition *minus-adelic-int-quot* ::

$'a \text{ adelic-int-quot} \Rightarrow 'a \text{ adelic-int-quot} \Rightarrow 'a \text{ adelic-int-quot}$

where $\text{minus-adelic-int-quot} \equiv (\lambda x y. x + - y)$

lemma *minus-adelic-int-quot-rep-eq*:

$\text{Rep-adelic-int-quot } (x - y) = \text{Rep-adelic-int-quot } x - \text{Rep-adelic-int-quot } y$

<proof>

lemma *minus-adelic-int-quot-abs-eq*:

$\text{Abs-adelic-int-quot } x - \text{Abs-adelic-int-quot } y = \text{Abs-adelic-int-quot } (x - y)$

<proof>

lemma *minus-adelic-int-quot*: $\text{adelic-int-quot } x - \text{adelic-int-quot } y = \text{adelic-int-quot } (x - y)$

<proof>

definition $x * y = \text{Abs-adelic-int-quot } (\text{Rep-adelic-int-quot } x * \text{Rep-adelic-int-quot } y)$

lemma *times-adelic-int-quot-rep-eq*:

$\text{Rep-adelic-int-quot } (x * y) = \text{Rep-adelic-int-quot } x * \text{Rep-adelic-int-quot } y$

<proof>

lemma *times-adelic-int-quot-abs-eq*:

$\text{Abs-adelic-int-quot } x * \text{Abs-adelic-int-quot } y = \text{Abs-adelic-int-quot } (x * y)$

<proof>

lemma *times-adelic-int-quot*: $\text{adelic-int-quot } x * \text{adelic-int-quot } y = \text{adelic-int-quot } (x * y)$

<proof>

instance

<proof>

end

instantiation *adelic-int-quot* ::

$(\text{nontrivial-factorial-unique-euclidean-bezout}) \{ \text{divide-trivial}, \text{inverse} \}$

begin

lift-definition *inverse-adelic-int-coset* ::
'a adelic-int coset-by-dist-sbgrp $\Rightarrow$ *'a adelic-int coset-by-dist-sbgrp*
is *inverse*
 $\langle$ *proof* $\rangle$

definition *inverse-adelic-int-quot* :: *'a adelic-int-quot* $\Rightarrow$ *'a adelic-int-quot*
where
inverse-adelic-int-quot $x \equiv$
Abs-adelic-int-quot (*inverse-adelic-int-coset* (*Rep-adelic-int-quot* x))

lemma *inverse-adelic-int-quot*: *inverse* (*adelic-int-quot* x) = *adelic-int-quot* (*inverse* x)
 $\langle$ *proof* $\rangle$

lemma *inverse-adelic-int-quot-0*[*simp*]: *inverse* ($0 :: 'a \text{ adelic-int-quot}$) = 0
 $\langle$ *proof* $\rangle$

lemma *inverse-adelic-int-quot-1*[*simp*]: *inverse* ($1 :: 'a \text{ adelic-int-quot}$) = 1
 $\langle$ *proof* $\rangle$

definition *divide-adelic-int-quot* ::
'a adelic-int-quot $\Rightarrow$ *'a adelic-int-quot* $\Rightarrow$ *'a adelic-int-quot*
where *divide-adelic-int-quot-def*[*simp*]: *divide-adelic-int-quot* $x \ y \equiv x * \text{inverse } y$

instance $\langle$ *proof* $\rangle$

end

lemma *divide-adelic-int-quot*: *adelic-int-quot* $x / \text{adelic-int-quot } y = \text{adelic-int-quot}$ (x / y)
for $x \ y :: 'a :: \text{nontrivial-factorial-unique-euclidean-bezout adelic-int}$
 $\langle$ *proof* $\rangle$

6.2 Equivalence relative to a prime

6.2.1 Equivalence

overloading

p-equal-adelic-int-coset $\equiv$
p-equal :: *'a :: nontrivial-factorial-idom prime* $\Rightarrow$
'a adelic-int coset-by-dist-sbgrp $\Rightarrow$
'a adelic-int coset-by-dist-sbgrp $\Rightarrow$ *bool*
p-equal-adelic-int-quot $\equiv$
p-equal :: *'a :: nontrivial-factorial-idom prime* $\Rightarrow$
'a adelic-int-quot $\Rightarrow$ *'a adelic-int-quot* $\Rightarrow$ *bool*
p-restrict-adelic-int-coset $\equiv$
p-restrict :: *'a adelic-int coset-by-dist-sbgrp* $\Rightarrow$
(*'a prime* $\Rightarrow$ *bool*) $\Rightarrow$ *'a adelic-int coset-by-dist-sbgrp*

$p\text{-restrict-adelic-int-quot} \equiv$
 $p\text{-restrict} :: 'a \text{ adelic-int-quot} \Rightarrow$
 $('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow 'a \text{ adelic-int-quot}$
begin

lift-definition $p\text{-equal-adelic-int-coset} ::$
 $'a :: \text{nontrivial-factorial-idom prime} \Rightarrow$
 $'a \text{ adelic-int coset-by-dist-sbgrp} \Rightarrow$
 $'a \text{ adelic-int coset-by-dist-sbgrp} \Rightarrow \text{bool}$
is $\lambda p \ x \ y. x \simeq_p y \vee ((x - y)^{\circ p}) \geq 1$
 $\langle \text{proof} \rangle$

definition $p\text{-equal-adelic-int-quot} ::$
 $'a :: \text{nontrivial-factorial-idom prime} \Rightarrow$
 $'a \text{ adelic-int-quot} \Rightarrow 'a \text{ adelic-int-quot} \Rightarrow \text{bool}$
where
 $p\text{-equal-adelic-int-quot } p \ x \ y \equiv$
 $(\text{Rep-adelic-int-quot } x) \simeq_p (\text{Rep-adelic-int-quot } y)$

lift-definition $p\text{-restrict-adelic-int-coset} ::$
 $'a :: \text{nontrivial-factorial-idom adelic-int coset-by-dist-sbgrp} \Rightarrow$
 $('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow 'a \text{ adelic-int coset-by-dist-sbgrp}$
is $\lambda x \ P. x \text{ prestrict } P$
 $\langle \text{proof} \rangle$

definition $p\text{-restrict-adelic-int-quot} ::$
 $'a :: \text{nontrivial-factorial-idom adelic-int-quot} \Rightarrow$
 $('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow 'a \text{ adelic-int-quot}$
where
 $p\text{-restrict-adelic-int-quot } x \ P \equiv$
 $\text{Abs-adelic-int-quot } ((\text{Rep-adelic-int-quot } x) \text{ prestrict } P)$

end

context
fixes $p :: 'a :: \text{nontrivial-factorial-idom prime}$
begin

lemma $p\text{-equal-adelic-int-coset-abs-eq0}:$
 $\text{distinguished-quotient-coset } x \simeq_p 0 \longleftrightarrow$
 $x \simeq_p 0 \vee (x^{\circ p}) \geq 1$
for $x :: 'a \text{ adelic-int}$
 $\langle \text{proof} \rangle$

lemma $p\text{-equal-adelic-int-coset-abs-eq1}:$
 $\text{distinguished-quotient-coset } x \simeq_p 1 \longleftrightarrow$
 $x \simeq_p 1 \vee ((x - 1)^{\circ p}) \geq 1$
for $x :: 'a \text{ adelic-int}$
 $\langle \text{proof} \rangle$

```

lemma p-equal-adelic-int-quot-abs-eq:
  (adelic-int-quot x)  $\simeq_p$  (adelic-int-quot y)  $\longleftrightarrow$ 
    x  $\simeq_p$  y  $\vee ((x - y)^{\circ p}) \geq 1$ 
for x y :: 'a adelic-int
  <proof>

lemma p-equal-adelic-int-quot-abs-eq0:
  adelic-int-quot x  $\simeq_p$  0  $\longleftrightarrow$ 
    x  $\simeq_p$  0  $\vee (x^{\circ p}) \geq 1$ 
for x :: 'a adelic-int
  <proof>

lemma p-equal-adelic-int-quot-abs-eq1:
  adelic-int-quot x  $\simeq_p$  1  $\longleftrightarrow$ 
    x  $\simeq_p$  1  $\vee ((x - 1)^{\circ p}) \geq 1$ 
for x :: 'a adelic-int
  <proof>

end

lemma p-restrict-adelic-int-quot-abs-eq:
  (adelic-int-quot x) prestrict P = adelic-int-quot (x prestrict P)
for P :: 'a::nontrivial-factorial-idom prime  $\Rightarrow$  bool and x :: 'a adelic-int
  <proof>

global-interpretation p-equality-adelic-int-quot:
  p-equality-no-zero-divisors-1
  p-equal :: 'a::nontrivial-euclidean-ring-cancel prime  $\Rightarrow$ 
    'a adelic-int-quot  $\Rightarrow$  'a adelic-int-quot  $\Rightarrow$  bool
  <proof>

overloading
  globally-p-equal-adelic-int-quot  $\equiv$ 
  globally-p-equal ::
    'a::nontrivial-euclidean-ring-cancel adelic-int-quot  $\Rightarrow$ 
    'a adelic-int-quot  $\Rightarrow$  bool

begin

definition globally-p-equal-adelic-int-quot ::
  'a::nontrivial-euclidean-ring-cancel adelic-int-quot  $\Rightarrow$ 
  'a adelic-int-quot  $\Rightarrow$  bool
where globally-p-equal-adelic-int-quot-def[simp]:
  globally-p-equal-adelic-int-quot = p-equality-adelic-int-quot.globally-p-equal

end

```

6.2.2 Embedding of constants

global-interpretation *global-p-equal-adelic-int-quot*:

global-p-equal-w-consts
 $p\text{-equal} :: 'a::\text{nontrivial-euclidean-ring-cancel prime} \Rightarrow$
 $'a \text{ adelic-int-quot} \Rightarrow 'a \text{ adelic-int-quot} \Rightarrow \text{bool}$
 $p\text{-restrict} ::$
 $'a \text{ adelic-int-quot} \Rightarrow ('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow$
 $'a \text{ adelic-int-quot}$
 $\lambda f. \text{ adelic-int-quot } (\text{adelic-int-consts } f)$
 $\langle \text{proof} \rangle$

abbreviation *adelic-int-quot-consts* $\equiv$ *global-p-equal-adelic-int-quot.consts*

abbreviation *adelic-int-quot-const* $\equiv$ *global-p-equal-adelic-int-quot.const*

lemma *p-equal-adelic-int-quot-consts-iff*:

$(\text{adelic-int-quot-consts } f) \simeq_p (\text{adelic-int-quot-consts } g) \longleftrightarrow$
 $(f \text{ } p) \text{ } p \bmod p = (g \text{ } p) \text{ } p \bmod p$
 $(\text{is } ?L \longleftrightarrow ?R)$
 $\langle \text{proof} \rangle$

lemma *p-equal0-adelic-int-quot-consts-iff*:

$(\text{adelic-int-quot-consts } f) \simeq_p 0 \longleftrightarrow (f \text{ } p) \text{ } p \bmod p = 0$
 $(\text{is } ?L \longleftrightarrow ?R)$
 $\langle \text{proof} \rangle$

lemma *adelic-int-quot-consts-eq-iff*:

$\text{adelic-int-quot-consts } f = \text{adelic-int-quot-consts } g \longleftrightarrow$
 $(\forall p. (f \text{ } p) \text{ } p \bmod p = (g \text{ } p) \text{ } p \bmod p)$
 $\langle \text{proof} \rangle$

lemma *adelic-int-quot-consts-eq0-iff*:

$(\text{adelic-int-quot-consts } f) = 0 \longleftrightarrow$
 $(\forall p. (f \text{ } p) \text{ } p \bmod p = 0)$
 $\langle \text{proof} \rangle$

lemmas

$p\text{-equal-adelic-int-quot-const-iff} =$
 $p\text{-equal-adelic-int-quot-consts-iff}[of \text{ } \lambda p. \text{ } - \text{ } \lambda p. \text{ } -]$
and $p\text{-equal0-adelic-int-quot-const-iff} = p\text{-equal0-adelic-int-quot-consts-iff}[of \text{ } \lambda p. \text{ } -]$
and $\text{adelic-int-quot-const-eq-iff} = \text{adelic-int-quot-consts-eq-iff}[of \text{ } \lambda p. \text{ } - \text{ } \lambda p. \text{ } -]$
and $\text{adelic-int-quot-const-eq0-iff} = \text{adelic-int-quot-consts-eq0-iff}[of \text{ } \lambda p. \text{ } -]$

lemma *adelic-int-quot-UNIV-eq-consts-range*:

$(\text{UNIV}::'a::\text{nontrivial-euclidean-ring-cancel adelic-int-quot set}) =$
 $\text{range } (\text{adelic-int-quot-consts})$
 $\langle \text{proof} \rangle$

lemma *adelic-int-quot-constsE*:

fixes $x :: 'a::\text{nontrivial-euclidean-ring-cancel adelic-int-quot}$
obtains f **where** $x = \text{adelic-int-quot-consts } f$
 $\langle \text{proof} \rangle$

lemma *adelic-int-quot-reduced-constsE*:
fixes $x :: 'a::\text{nontrivial-euclidean-ring-cancel adelic-int-quot}$
obtains f
where $x = \text{adelic-int-quot-consts } f$
and $\forall p::'a \text{ prime. } (f\ p) \text{ pdiv } p = 0$
 $\langle \text{proof} \rangle$

lemma *adelic-int-quot-prestrict-zero-depth-abs-eq*:
 $\text{adelic-int-quot } (x \text{ prestrict } (\lambda p::'a \text{ prime. } (x^{\circ p}) = 0)) =$
 $\text{adelic-int-quot } x$
for $x :: 'a::\text{nontrivial-euclidean-ring-cancel adelic-int}$
 $\langle \text{proof} \rangle$

6.2.3 Division and inverses

global-interpretation *p-equal-div-inv-adelic-int-quot*:
 $\text{global-p-equal-div-inv}$
 $p\text{-equal} :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow$
 $'a \text{ adelic-int-quot} \Rightarrow 'a \text{ adelic-int-quot} \Rightarrow \text{bool}$
 $p\text{-restrict} ::$
 $'a \text{ adelic-int-quot} \Rightarrow ('a \text{ prime} \Rightarrow \text{bool}) \Rightarrow$
 $'a \text{ adelic-int-quot}$
 $\langle \text{proof} \rangle$

6.3 Special case of integer

context
fixes $p :: \text{int prime}$
begin

lemma *inj-on-const-prestrict-single-int-prime*:
 $\text{inj-on } (\lambda k. \text{ adelic-int-quot-const } k \text{ prestrict } ((=) p)) \{0..\text{Rep-prime } p - 1\}$
 $\langle \text{proof} \rangle$

lemma *range-prestrict-single-int-prime*:
 $\text{range } (\lambda x. x \text{ prestrict } ((=) p)) =$
 $(\lambda k. \text{ adelic-int-quot-const } k \text{ prestrict } ((=) p)) \text{ ` } \{0..\text{Rep-prime } p - 1\}$
 $\langle \text{proof} \rangle$

lemma *finite-range-prestrict-single-int-prime*:
 $\text{finite } (\text{range } (\lambda x::\text{int adelic-int-quot. } x \text{ prestrict } ((=) p)))$
 $\langle \text{proof} \rangle$

lemma *card-range-prestrict-single-int-prime*:
 $\text{card } (\text{range } (\lambda x::\text{int adelic-int-quot. } x \text{ prestrict } ((=) p))) = \text{nat } (\text{Rep-prime } p)$
 $\langle \text{proof} \rangle$

end

end

theory *More-pAdic-Product*

imports *Fin-Field-Product*

begin

7 Compactness of local ring of integers

7.1 Preliminaries

lemma *infinite-vimageE*:

fixes $f :: 'a \Rightarrow 'b$

assumes *infinite* ($UNIV :: 'a \text{ set}$) **and** $\text{range } f \subseteq B$ **and** *finite* B

obtains b **where** $b \in B$ **and** *infinite* ($f^{-1}\{b\}$)

<proof>

primrec *subset-subseq* ::

$(\text{nat} \Rightarrow 'a) \Rightarrow 'a \text{ set} \Rightarrow \text{nat} \Rightarrow \text{nat}$

where

$\text{subset-subseq } f \ A \ 0 = (\text{LEAST } k. f \ k \in A) \mid$

$\text{subset-subseq } f \ A \ (\text{Suc } n) = (\text{LEAST } k. k > \text{subset-subseq } f \ A \ n \wedge f \ k \in A)$

lemma

assumes *infinite* ($f^{-1} A$)

shows *subset-subseq-strict-mono*: *strict-mono* ($\text{subset-subseq } f \ A$)

and *subset-subseq-range* : $f (\text{subset-subseq } f \ A \ n) \in A$

<proof>

7.2 Sequential compactness

7.2.1 Creating a Cauchy subsequence

abbreviation

$p\text{-adic-prod-int-convergent-subseq-seq-step } p \ X \ g \ n \equiv$

$\lambda k. p\text{-adic-prod-int-quot } (p\text{-adic-prod-shift-}p\text{-depth } p \ (- \text{int } n) ((X \ (g \ k) - X \ (g \ 0))))$

primrec *p-adic-prod-int-convergent-subseq-seq* ::

$'a::\text{nontrivial-factorial-unique-euclidean-bezout prime} \Rightarrow$

$(\text{nat} \Rightarrow 'a \text{ } p\text{-adic-prod}) \Rightarrow \text{nat} \Rightarrow (\text{nat} \Rightarrow \text{nat})$

where *p-adic-prod-int-convergent-subseq-seq* $p \ X \ 0 = \text{id} \mid$

p-adic-prod-int-convergent-subseq-seq $p \ X \ (\text{Suc } n) =$

p-adic-prod-int-convergent-subseq-seq $p \ X \ n \circ$

subset-subseq

```

(
  p-adic-prod-int-convergent-subseq-seq-step p X (
    p-adic-prod-int-convergent-subseq-seq p X n
  ) n
)
{SOME z.
  infinite (
    p-adic-prod-int-convergent-subseq-seq-step p X (
      p-adic-prod-int-convergent-subseq-seq p X n
    ) n
    - ' {z}
  )
}

```

abbreviation

```

p-adic-prod-int-convergent-subseq p X n ≡
p-adic-prod-int-convergent-subseq-seq p X n n

```

lemma restricted-X-infinite-quot-vimage:

```

∃ z. infinite ((λk. p-adic-prod-int-quot (X k)) - ' {z})
if fin-p-quot : finite (range (λx::'a adelic-int-quot. x prestrict ((=) p)))
and range-X : range X ⊆ O∇p
and restricted-X: ∀ n. X n = (X n) prestrict ((=) p)
for p :: 'a::nontrivial-factorial-idom prime and X :: nat ⇒ 'a p-adic-prod
⟨proof⟩

```

context

```

fixes p      :: 'a::nontrivial-factorial-unique-euclidean-bezout prime
and X        :: nat ⇒ 'a p-adic-prod
and ss       :: nat ⇒ (nat ⇒ nat)
and ss-step  :: nat ⇒ (nat ⇒ 'a adelic-int-quot)
and nth-im   :: nat ⇒ 'a adelic-int-quot
and subss    :: nat ⇒ (nat ⇒ nat)
defines
  ss ≡ p-adic-prod-int-convergent-subseq-seq p X
and
  ss-step ≡ λn. p-adic-prod-int-convergent-subseq-seq-step p X (ss n) n
and nth-im ≡ λn. (SOME z. infinite (ss-step n - ' {z}))
and subss ≡ λn. subset-subseq (ss-step n) {nth-im n}
assumes fin-p-quot : finite (range (λx::'a adelic-int-quot. x prestrict ((=) p)))
and range-X : range X ⊆ O∇p
and restricted-X: ∀ n. X n = (X n) prestrict ((=) p)
begin

```

lemma p-adic-prod-int-convergent-subseq-seq-step-infinite-quot-vimage:

```

∃ z. infinite (ss-step n - ' {z}) (is ?A)
and p-adic-prod-int-convergent-subseq-seq-step-depth:
  ∀ k. (X (ss n k)) ⊆p (X (ss n 0)) →
    ((X (ss n k) - X (ss n 0))op) ≥ int n

```

(is ?B)
 <proof>

lemma *p-adic-prod-int-convergent-subset-subseq-strict-mono*:
 strict-mono (subset-subseq (ss-step n) {nth-im n})
 <proof>

lemma *p-adic-prod-int-convergent-subseq-strict-mono'*: strict-mono (ss n)
 <proof>

lemma *p-adic-prod-int-convergent-subseq-strict-mono*:
 strict-mono (p-adic-prod-int-convergent-subseq p X)
 <proof>

lemma *p-adic-prod-int-convergent-subseq-Cauchy*:
 p-adic-prod-p-cauchy p (X $\circ$ p-adic-prod-int-convergent-subseq p X)
 <proof>

lemma *p-adic-prod-int-convergent-subseq-limseq*:
 $x \in (\lambda z. z \text{ prestrict } ((=) p)) \text{ ' } \mathcal{O}_{\forall p}$
 if p-adic-prod-p-open-nbhds-limseq (X $\circ$ g) x
 <proof>

end

7.2.2 Proving sequential compactness

context
 fixes p :: 'a::nontrivial-factorial-unique-euclidean-bezout prime
 assumes fin-p-quot: finite (range ($\lambda x::'a$ adelic-int-quot. x prestrict ((=) p)))
 begin

lemma *p-adic-prod-local-int-ring-seq-compact'*:
 $\exists g::nat \Rightarrow nat. \text{ strict-mono } g \wedge \text{ p-adic-prod-p-cauchy } p (X \circ g)$
 if range-X:
 $\text{ range } X \subseteq (\lambda x. x \text{ prestrict } ((=) p)) \text{ ' } (\mathcal{O}_{\forall p})$
 for X :: nat $\Rightarrow$ 'a p-adic-prod
 <proof>

lemma *p-adic-prod-local-int-ring-seq-compact*:
 $\exists x \in (\lambda x. x \text{ prestrict } ((=) p)) \text{ ' } (\mathcal{O}_{\forall p}).$
 $\exists g::nat \Rightarrow nat.$
 $\text{ strict-mono } g \wedge \text{ p-adic-prod-p-open-nbhds-limseq } (X \circ g) x$
 if range-X:
 $\text{ range } X \subseteq (\lambda x. x \text{ prestrict } ((=) p)) \text{ ' } (\mathcal{O}_{\forall p})$
 for X :: nat $\Rightarrow$ 'a p-adic-prod
 <proof>

lemma *adelic-int-local-int-ring-seq-abs*:

```

range X ⊆ (λx. x prestrict ((=) p)) ‘ (O∇p)
if range-X: range X ⊆ O∇p
and range-abs-X:
  range (λk. Abs-adelic-int (X k)) ⊆ range (λx. x prestrict ((=) p))
for X :: nat ⇒ 'a p-adic-prod
⟨proof⟩

```

```

lemma adelic-int-local-int-ring-seq-compact':
  ∃ g::nat⇒nat. strict-mono g ∧ adelic-int-p-cauchy p (X ∘ g)
if range-X:
  range X ⊆ (λx. x prestrict ((=) p)) ‘ (O∇p)
for X :: nat ⇒ 'a adelic-int
⟨proof⟩

```

```

lemma adelic-int-local-int-ring-seq-compact:
  ∃ x∈range (λx. x prestrict ((=) p)).
  ∃ g::nat⇒nat.
    strict-mono g ∧ adelic-int-p-open-nbhds-limseq (X ∘ g) x
if range-X:
  range X ⊆ range (λx. x prestrict ((=) p))
for X :: nat ⇒ 'a adelic-int
⟨proof⟩

```

end

```

lemma int-adic-prod-local-int-ring-seq-compact:
  ∃ x∈(λx. x prestrict ((=) p)) ‘ (O∇p).
  ∃ g::nat⇒nat.
    strict-mono g ∧ p-adic-prod-p-open-nbhds-limseq (X ∘ g) x
if range X ⊆ (λx. x prestrict ((=) p)) ‘ (O∇p)
for p :: int prime
and X :: nat ⇒ int p-adic-prod
⟨proof⟩

```

```

lemma int-adelic-int-local-int-ring-seq-compact:
  ∃ x∈range (λx. x prestrict ((=) p)).
  ∃ g::nat⇒nat.
    strict-mono g ∧ adelic-int-p-open-nbhds-limseq (X ∘ g) x
if range X ⊆ range (λx. x prestrict ((=) p))
for p :: int prime
and X :: nat ⇒ int adelic-int
⟨proof⟩

```

7.3 Finite-open-cover compactness

```

function exclusion-sequence ::
  'a set ⇒ ('a ⇒ 'a set) ⇒
  'a ⇒ nat ⇒ 'a
where exclusion-sequence A f a 0 = a |

```

$\text{exclusion-sequence } A \text{ f } a \text{ (Suc } n) =$
 $(\text{SOME } x. x \in A \wedge x \notin (\bigcup k \leq n. f \text{ (exclusion-sequence } A \text{ f } a \text{ k)}))$
 $\langle \text{proof} \rangle$
termination $\langle \text{proof} \rangle$

lemma
assumes $\neg A \subseteq (\bigcup k \leq n. f \text{ (exclusion-sequence } A \text{ f } a \text{ k)})$
shows $\text{exclusion-sequence-mem: exclusion-sequence } A \text{ f } a \text{ (Suc } n) \in A$
and $\text{exclusion-sequence-excludes:}$
 $\text{exclusion-sequence } A \text{ f } a \text{ (Suc } n) \notin (\bigcup k \leq n. f \text{ (exclusion-sequence } A \text{ f } a \text{ k)})$
 $\langle \text{proof} \rangle$

context
fixes $p :: 'a::\text{nontrivial-factorial-unique-euclidean-bezout prime}$
assumes $\text{fin-p-quot: finite (range } (\lambda x::'a \text{ adelic-int-quot. } x \text{ prestrict } ((=) \text{ p})))$
begin

lemma $p\text{-adic-prod-local-int-ring-finite-cover:}$
 $\exists C. \text{finite } C \wedge$
 $C \subseteq (\lambda x. x \text{ prestrict } ((=) \text{ p})) ' \mathcal{O}_{\forall p} \wedge$
 $(\lambda x. x \text{ prestrict } ((=) \text{ p})) ' \mathcal{O}_{\forall p} \subseteq$
 $(\bigcup c \in C. p\text{-adic-prod-p-open-nbhd } p \text{ n } c)$
for $C :: 'a \text{ set}$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-local-int-ring-lebesgue-number:}$
 $\exists d. \forall x \in (\lambda x. x \text{ prestrict } ((=) \text{ p})) ' \mathcal{O}_{\forall p}.$
 $\exists A \in \mathcal{A}. p\text{-adic-prod-p-open-nbhd } p \text{ d } x \subseteq A$
if cover:
 $(\lambda x. x \text{ prestrict } ((=) \text{ p})) ' \mathcal{O}_{\forall p} \subseteq \bigcup \mathcal{A}$
and by-opens:
 $\forall A \in \mathcal{A}. \text{generate-topology } (p\text{-adic-prod-local-p-open-nbhds } p) A$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-local-int-ring-compact:}$
 $\exists \mathcal{B} \subseteq \mathcal{A}. \text{finite } \mathcal{B} \wedge$
 $(\lambda x. x \text{ prestrict } ((=) \text{ p})) ' \mathcal{O}_{\forall p} \subseteq \bigcup \mathcal{B}$
if cover:
 $(\lambda x. x \text{ prestrict } ((=) \text{ p})) ' \mathcal{O}_{\forall p} \subseteq \bigcup \mathcal{A}$
and by-opens:
 $\forall A \in \mathcal{A}. \text{generate-topology } (p\text{-adic-prod-local-p-open-nbhds } p) A$
 $\langle \text{proof} \rangle$

lemma $p\text{-adic-prod-local-scaled-int-ring-compact:}$
 $\exists \mathcal{B} \subseteq \mathcal{A}. \text{finite } \mathcal{B} \wedge$
 $(\lambda x. x \text{ prestrict } ((=) \text{ p})) ' (\mathcal{P}_{\forall p}^n)$
 $\subseteq \bigcup \mathcal{B}$
if cover:
 $(\lambda x. x \text{ prestrict } ((=) \text{ p})) ' (\mathcal{P}_{\forall p}^n)$

$\subseteq \bigcup \mathcal{A}$
and *by-opens*:
 $\forall A \in \mathcal{A}. \text{generate-topology } (p\text{-adic-prod-local-}p\text{-open-nbhds } p) A$
for $\mathcal{A} :: 'a \text{ } p\text{-adic-prod set set}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-local-int-ring-compact*:
 $\exists \mathcal{B} \subseteq \mathcal{A}. \text{finite } \mathcal{B} \wedge$
 $\text{range } (\lambda x. x \text{ prestrict } ((=) p)) \subseteq \bigcup \mathcal{B}$
if *fin-p-quot*: $\text{finite } (\text{range } (\lambda x :: 'a \text{ adelic-int-quot}. x \text{ prestrict } ((=) p)))$
and *cover*: $\text{range } (\lambda x. x \text{ prestrict } ((=) p)) \subseteq \bigcup \mathcal{A}$
and *by-opens*:
 $\forall A \in \mathcal{A}. \text{generate-topology } (\text{adelic-int-local-}p\text{-open-nbhds } p) A$
for $\mathcal{A} :: 'a \text{ adelic-int set set}$
 $\langle \text{proof} \rangle$

lemma *adelic-int-local-scaled-int-ring-compact*:
 $\exists \mathcal{B} \subseteq \mathcal{A}. \text{finite } \mathcal{B} \wedge$
 $(\lambda x. x \text{ prestrict } ((=) p)) ' (\mathcal{P}_{\forall p}^n)$
 $\subseteq \bigcup \mathcal{B}$
if *fin-p-quot*: $\text{finite } (\text{range } (\lambda x :: 'a \text{ adelic-int-quot}. x \text{ prestrict } ((=) p)))$
and *cover*:
 $(\lambda x. x \text{ prestrict } ((=) p)) ' (\mathcal{P}_{\forall p}^n)$
 $\subseteq \bigcup \mathcal{A}$
and *by-opens*:
 $\forall A \in \mathcal{A}. \text{generate-topology } (\text{adelic-int-local-}p\text{-open-nbhds } p) A$
for $\mathcal{A} :: 'a \text{ adelic-int set set}$
 $\langle \text{proof} \rangle$

end

lemma *int-adic-prod-local-int-ring-compact*:
 $\exists \mathcal{B} \subseteq \mathcal{A}. \text{finite } \mathcal{B} \wedge$
 $(\lambda x. x \text{ prestrict } ((=) p)) ' \mathcal{O}_{\forall p} \subseteq \bigcup \mathcal{B}$
if $(\lambda x. x \text{ prestrict } ((=) p)) ' \mathcal{O}_{\forall p} \subseteq \bigcup \mathcal{A}$
and $\forall A \in \mathcal{A}. \text{generate-topology } (p\text{-adic-prod-local-}p\text{-open-nbhds } p) A$
for $p :: \text{int prime}$
and $\mathcal{A} :: \text{int } p\text{-adic-prod set set}$
 $\langle \text{proof} \rangle$

lemma *int-adic-prod-local-scaled-int-ring-compact*:
 $\exists \mathcal{B} \subseteq \mathcal{A}. \text{finite } \mathcal{B} \wedge$
 $(\lambda x. x \text{ prestrict } ((=) p)) ' (\mathcal{P}_{\forall p}^n)$
 $\subseteq \bigcup \mathcal{B}$
if
 $(\lambda x. x \text{ prestrict } ((=) p)) ' (\mathcal{P}_{\forall p}^n)$
 $\subseteq \bigcup \mathcal{A}$
and $\forall A \in \mathcal{A}. \text{generate-topology } (p\text{-adic-prod-local-}p\text{-open-nbhds } p) A$
for $p :: \text{int prime}$

```

and  $\mathcal{A} :: \text{int } p\text{-adic-prod set set}$ 
 $\langle \text{proof} \rangle$ 

lemma int-adelic-int-local-int-ring-compact:
   $\exists \mathcal{B} \subseteq \mathcal{A}. \text{finite } \mathcal{B} \wedge$ 
     $\text{range } (\lambda x. x \text{ prestrict } ((=) p)) \subseteq \bigcup \mathcal{B}$ 
  if  $\text{range } (\lambda x. x \text{ prestrict } ((=) p)) \subseteq \bigcup \mathcal{A}$ 
  and  $\forall A \in \mathcal{A}. \text{generate-topology } (\text{adelic-int-local-}p\text{-open-nbhd } p) A$ 
  for  $p :: \text{int prime}$ 
  and  $\mathcal{A} :: \text{int adelic-int set set}$ 
   $\langle \text{proof} \rangle$ 

lemma int-adelic-int-local-scaled-int-ring-compact:
   $\exists \mathcal{B} \subseteq \mathcal{A}. \text{finite } \mathcal{B} \wedge$ 
     $(\lambda x. x \text{ prestrict } ((=) p)) ' (\mathcal{P}_{\forall p}^n)$ 
     $\subseteq \bigcup \mathcal{B}$ 
  if
     $(\lambda x. x \text{ prestrict } ((=) p)) ' (\mathcal{P}_{\forall p}^n)$ 
     $\subseteq \bigcup \mathcal{A}$ 
  and  $\forall A \in \mathcal{A}. \text{generate-topology } (\text{adelic-int-local-}p\text{-open-nbhd } p) A$ 
  for  $p :: \text{int prime}$ 
  and  $\mathcal{A} :: \text{int adelic-int set set}$ 
   $\langle \text{proof} \rangle$ 

end

```

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