

Complex Bounded Operators*

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Abstract

We present a formalization of bounded operators on complex vector spaces. Our formalization contains material on complex vector spaces (normed spaces, Banach spaces, Hilbert spaces) that complements and goes beyond the developments of real vectors spaces in the Isabelle/HOL standard library. We define the type of bounded operators between complex vector spaces (*cblinfun*) and develop the theory of unitaries, projectors, extension of bounded linear functions (BLT theorem), adjoints, Loewner order, closed subspaces and more. For the finite-dimensional case, we provide code generation support by identifying finite-dimensional operators with matrices as formalized in the *Jordan_Normal_Form* AFP entry.

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Theories whose names end with 0 are complex analogues of the similarly named theories concerning real vector spaces in the Isabelle/HOL standard library. They are kept in sync with their real counterparts. The theories without 0 contain material that goes beyond the material in the Isabelle/HOL standard library. This separation allows to keep the material in sync more easily when the Isabelle/HOL standard library is updated.

1 *Extra-Pretty-Code-Examples* – Setup for nicer output of *value*

```

theory Extra-Pretty-Code-Examples
imports
  HOL-Examples.Sqrt
  Real-Impl.Real-Impl
  HOL-Library.Code-Target-Numeral
  Jordan-Normal-Form.Matrix-Impl
begin

```

Some setup that makes the output of the *value* command more readable if matrices and complex numbers are involved.

It is not recommended to import this theory in theories that get included in actual developments (because of the changes to the code generation setup).

It is meant for inclusion in example theories only.

```

lemma two-sqrt-irrat[simp]:  $2 \in \text{sqrt-irrat}$ 
  <proof>

```

```

lemma complex-number-code-post[code-post]:
  shows Complex  $a\ 0 = \text{complex-of-real } a$ 
  and complex-of-real  $0 = 0$ 

```

```

and complex-of-real 1 = 1
and complex-of-real (a/b) = complex-of-real a / complex-of-real b
and complex-of-real (numeral n) = numeral n
and complex-of-real (-r) = - complex-of-real r
⟨proof⟩

```

```

lemma real-number-code-post[code-post]:
shows real-of (Abs-mini-alg (p, 0, b)) = real-of-rat p
and real-of (Abs-mini-alg (p, q, 2)) = real-of-rat p + real-of-rat q * sqrt 2
and sqrt 0 = 0
and sqrt (real 0) = 0
and x * (0::real) = 0
and (0::real) * x = 0
and (0::real) + x = x
and x + (0::real) = x
and (1::real) * x = x
and x * (1::real) = x
⟨proof⟩

```

```

translations x ← CONST IArray x

```

```

end

```

2 Extra-General – General missing things

```

theory Extra-General
imports
  HOL-Library.Cardinality
  HOL-Analysis.Elementary-Topology
  HOL-Analysis.Uniform-Limit
  HOL-Library.Set-Algebras
  HOL-Types-To-Sets.Types-To-Sets
  HOL-Library.Complex-Order
  HOL-Analysis.Infinite-Sum
  HOL-Cardinals.Cardinals
  HOL-Library.Complemented-Lattices
  HOL-Analysis.Abstract-Topological-Spaces
begin

```

2.1 Misc

```

lemma reals-zero-comparable:
fixes x::complex
assumes x∈ℝ
shows x ≤ 0 ∨ x ≥ 0
⟨proof⟩

```

lemma *unique-choice*: $\forall x. \exists! y. Q\ x\ y \implies \exists! f. \forall x. Q\ x\ (f\ x)$
 $\langle proof \rangle$

lemma *image-set-plus*:
assumes $\langle linear\ U \rangle$
shows $\langle U\ ' (A + B) = U\ ' A + U\ ' B \rangle$
 $\langle proof \rangle$

consts *heterogenous-identity* :: $\langle 'a \Rightarrow 'b \rangle$
overloading *heterogenous-identity-id* $\equiv$ *heterogenous-identity* :: $'a \Rightarrow 'a$ **begin**
definition *heterogenous-identity-def[simp]*: $\langle heterogenous-identity-id = id \rangle$
end

lemma *L2-set-mono2*:
assumes *a1*: *finite* *L* **and** *a2*: $K \leq L$
shows $L2\text{-set}\ f\ K \leq L2\text{-set}\ f\ L$
 $\langle proof \rangle$

lemma *Sup-real-close*:
fixes *e* :: *real*
assumes $0 < e$
and *S*: *bdd-above* *S* $S \neq \{\}$
shows $\exists x \in S. Sup\ S - e < x$
 $\langle proof \rangle$

Improved version of *internalize-sort*: It is not necessary to specify the sort of the type variable.

$\langle ML \rangle$

lemma *card-prod-omega*: $\langle X *c\ natLeq =o\ X \rangle$ **if** $\langle C\infinite\ X \rangle$
 $\langle proof \rangle$

lemma *countable-leq-natLeq*: $\langle |X| \leq_o\ natLeq \rangle$ **if** $\langle countable\ X \rangle$
 $\langle proof \rangle$

lemma *set-Times-plus-distrib*: $\langle (A \times B) + (C \times D) = (A + C) \times (B + D) \rangle$
 $\langle proof \rangle$

definition (**in order**) $\langle is\text{-}Sup\ X\ s \longleftrightarrow (\forall x \in X. x \leq s) \wedge (\forall y. (\forall x \in X. x \leq y) \longrightarrow s \leq y) \rangle$

definition (**in order**) $\langle has\text{-}Sup\ X \longleftrightarrow (\exists s. is\text{-}Sup\ X\ s) \rangle$

lemma (**in order**) *is-SupI*:
assumes $\langle \bigwedge x. x \in X \implies x \leq s \rangle$
assumes $\langle \bigwedge y. (\bigwedge x. x \in X \implies x \leq y) \implies s \leq y \rangle$
shows $\langle is\text{-}Sup\ X\ s \rangle$
 $\langle proof \rangle$

```

lemma is-Sup-approx-below:
  fixes  $b :: \langle 'a::\text{linordered-ab-group-add} \rangle$ 
  assumes  $\langle \text{is-Sup } A \ b \rangle$ 
  assumes  $\langle \varepsilon > 0 \rangle$ 
  shows  $\langle \exists x \in A. \ b - \varepsilon \leq x \rangle$ 
 $\langle \text{proof} \rangle$ 

```

2.2 Not singleton

```

class not-singleton =
  assumes not-singleton-card:  $\exists x \ y. \ x \neq y$ 

```

```

lemma not-singleton-existence[simp]:
   $\langle \exists x::('a::\text{not-singleton}). \ x \neq t \rangle$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma not-not-singleton-zero:
   $\langle x = 0 \rangle$  if  $\langle \neg \text{class.not-singleton TYPE}('a) \rangle$  for  $x :: \langle 'a::\text{zero} \rangle$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma UNIV-not-singleton[simp]:  $(\text{UNIV}:::\text{not-singleton set}) \neq \{x\}$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma UNIV-not-singleton-converse:
  assumes  $\bigwedge x::'a. \ \text{UNIV} \neq \{x\}$ 
  shows  $\exists x::'a. \ \exists y. \ x \neq y$ 
 $\langle \text{proof} \rangle$ 

```

```

subclass (in card2) not-singleton
 $\langle \text{proof} \rangle$ 

```

```

subclass (in perfect-space) not-singleton
 $\langle \text{proof} \rangle$ 

```

```

lemma class-not-singletonI-monoid-add:
  assumes  $(\text{UNIV}::'a \text{ set}) \neq \{0\}$ 
  shows  $\text{class.not-singleton TYPE}('a::\text{monoid-add})$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma not-singleton-vs-CARD-1:
  assumes  $\langle \neg \text{class.not-singleton TYPE}('a) \rangle$ 
  shows  $\langle \text{class.CARD-1 TYPE}('a) \rangle$ 
 $\langle \text{proof} \rangle$ 

```

2.3 CARD-1

```

context CARD-1 begin

```

```

lemma everything-the-same[simp]:  $(x::'a)=y$ 
 $\langle \text{proof} \rangle$ 

```


lemma *CARD-1-UNIV*: $UNIV = \{x::'a\}$
 ⟨proof⟩

lemma *CARD-1-ext*: $x (a::'a) = y \ b \implies x = y$
 ⟨proof⟩

end

instance *unit* :: *CARD-1*
 ⟨proof⟩

instance *prod* :: (*CARD-1*, *CARD-1*) *CARD-1*
 ⟨proof⟩

instance *fun* :: (*CARD-1*, *CARD-1*) *CARD-1*
 ⟨proof⟩

lemma *enum-CARD-1*: $(Enum.enum :: 'a::\{CARD-1,enum\} list) = [a]$
 ⟨proof⟩

lemma *card-not-singleton*: $\langle CARD('a::not-singleton) \neq 1 \rangle$
 ⟨proof⟩

2.4 Topology

lemma *cauchy-filter-metricI*:
fixes $F :: 'a::metric-space filter$
assumes $\bigwedge e. e > 0 \implies \exists P. eventually\ P\ F \wedge (\forall x\ y. P\ x \wedge P\ y \longrightarrow dist\ x\ y < e)$
shows *cauchy-filter* F
 ⟨proof⟩

lemma *cauchy-filter-metric-filtermapI*:
fixes $F :: 'a filter$ **and** $f :: 'a \Rightarrow 'b::metric-space$
assumes $\bigwedge e. e > 0 \implies \exists P. eventually\ P\ F \wedge (\forall x\ y. P\ x \wedge P\ y \longrightarrow dist\ (f\ x)\ (f\ y) < e)$
shows *cauchy-filter* $(filtermap\ f\ F)$
 ⟨proof⟩

lemma *tendsto-add-const-iff*:
 — This is a generalization of *Limits.tendsto-add-const-iff*, the only difference is that the sort here is more general.
 $((\lambda x. c + f\ x :: 'a::topological-group-add) \longrightarrow c + d)\ F \longleftrightarrow (f \longrightarrow d)\ F$
 ⟨proof⟩

lemma *finite-subsets-at-top-minus*:

assumes $A \subseteq B$
shows $\text{finite-subsets-at-top } (B - A) \leq \text{filtermap } (\lambda F. F - A) (\text{finite-subsets-at-top } B)$
 $\langle \text{proof} \rangle$

lemma *finite-subsets-at-top-inter*:
assumes $A \subseteq B$
shows $\text{filtermap } (\lambda F. F \cap A) (\text{finite-subsets-at-top } B) = \text{finite-subsets-at-top } A$
 $\langle \text{proof} \rangle$

lemma *tendsto-principal-singleton*:
shows $(f \longrightarrow f\ x) (\text{principal } \{x\})$
 $\langle \text{proof} \rangle$

lemma *complete-singleton*:
 $\text{complete } \{s :: 'a :: \text{uniform-space}\}$
 $\langle \text{proof} \rangle$

lemma *on-closure-eqI*:
fixes $f\ g :: \langle 'a :: \text{topological-space} \Rightarrow 'b :: t2\text{-space} \rangle$
assumes $\text{eq}: \langle \bigwedge x. x \in S \implies f\ x = g\ x \rangle$
assumes $xS: \langle x \in \text{closure } S \rangle$
assumes $\text{cont}: \langle \text{continuous-on } \text{UNIV } f \rangle \langle \text{continuous-on } \text{UNIV } g \rangle$
shows $\langle f\ x = g\ x \rangle$
 $\langle \text{proof} \rangle$

lemma *on-closure-leI*:
fixes $f\ g :: \langle 'a :: \text{topological-space} \Rightarrow 'b :: \text{linorder-topology} \rangle$
assumes $\text{eq}: \langle \bigwedge x. x \in S \implies f\ x \leq g\ x \rangle$
assumes $xS: \langle x \in \text{closure } S \rangle$
assumes $\text{cont}: \langle \text{continuous-on } \text{UNIV } f \rangle \langle \text{continuous-on } \text{UNIV } g \rangle$
shows $\langle f\ x \leq g\ x \rangle$
 $\langle \text{proof} \rangle$

lemma *tendsto-compose-at-within*:
assumes $f: (f \longrightarrow y) F$ **and** $g: (g \longrightarrow z) (\text{at } y \text{ within } S)$
and $\text{fg}: \text{eventually } (\lambda w. f\ w = y \longrightarrow g\ y = z) F$
and $fS: \langle \forall_F w \text{ in } F. f\ w \in S \rangle$
shows $((g \circ f) \longrightarrow z) F$
 $\langle \text{proof} \rangle$

2.5 Sums

lemma *sum-single*:
assumes *finite* A
assumes $\bigwedge j. j \neq i \implies j \in A \implies f\ j = 0$
shows $\text{sum } f\ A = (\text{if } i \in A \text{ then } f\ i \text{ else } 0)$
 $\langle \text{proof} \rangle$

lemma *has-sum-comm-additive-general*:

— This is a strengthening of *has-sum-comm-additive-general*.

fixes $f :: \langle 'b :: \{ \text{comm-monoid-add}, \text{topological-space} \} \Rightarrow 'c :: \{ \text{comm-monoid-add}, \text{topological-space} \} \rangle$

assumes $f\text{-sum}: \langle \bigwedge F. \text{finite } F \Longrightarrow F \subseteq S \Longrightarrow \text{sum } (f \circ g) F = f (\text{sum } g F) \rangle$

— Not using *additive* because it would add sort constraint *ab-group-add*

assumes $\text{inS}: \langle \bigwedge F. \text{finite } F \Longrightarrow \text{sum } g F \in T \rangle$

assumes $\text{cont}: \langle (f \longrightarrow f x) \text{ (at } x \text{ within } T) \rangle$

— For *t2-space* and $T = \text{UNIV}$, this is equivalent to *isCont* $f x$ by *isCont-def*.

assumes $\text{infsum}: \langle (g \text{ has-sum } x) S \rangle$

shows $\langle ((f \circ g) \text{ has-sum } (f x)) S \rangle$

<proof>

lemma *summable-on-comm-additive-general*:

— This is a strengthening of *summable-on-comm-additive-general*.

fixes $g :: \langle 'a \Rightarrow 'b :: \{ \text{comm-monoid-add}, \text{topological-space} \} \rangle$ **and** $f :: \langle 'b \Rightarrow 'c :: \{ \text{comm-monoid-add}, \text{topological-space} \} \rangle$

assumes $\langle \bigwedge F. \text{finite } F \Longrightarrow F \subseteq S \Longrightarrow \text{sum } (f \circ g) F = f (\text{sum } g F) \rangle$

— Not using *additive* because it would add sort constraint *ab-group-add*

assumes $\text{inS}: \langle \bigwedge F. \text{finite } F \Longrightarrow \text{sum } g F \in T \rangle$

assumes $\text{cont}: \langle \bigwedge x. (g \text{ has-sum } x) S \Longrightarrow (f \longrightarrow f x) \text{ (at } x \text{ within } T) \rangle$

— For *t2-space* and $T = \text{UNIV}$, this is equivalent to *isCont* $f x$ by *isCont-def*.

assumes $\langle g \text{ summable-on } S \rangle$

shows $\langle (f \circ g) \text{ summable-on } S \rangle$

<proof>

lemma *has-sum-metric*:

fixes $l :: \langle 'a :: \{ \text{metric-space}, \text{comm-monoid-add} \} \rangle$

shows $\langle (f \text{ has-sum } l) A \longleftrightarrow (\forall e. e > 0 \longrightarrow (\exists X. \text{finite } X \wedge X \subseteq A \wedge (\forall Y. \text{finite } Y \wedge X \subseteq Y \wedge Y \subseteq A \longrightarrow \text{dist } (\text{sum } f Y) l < e))) \rangle$

<proof>

lemma *summable-on-product-finite-left*:

fixes $f :: \langle 'a \times 'b \Rightarrow 'c :: \{ \text{topological-comm-monoid-add} \} \rangle$

assumes $\text{sum}: \langle \bigwedge x. x \in X \Longrightarrow (\lambda y. f(x, y)) \text{ summable-on } Y \rangle$

assumes $\langle \text{finite } X \rangle$

shows $\langle f \text{ summable-on } (X \times Y) \rangle$

<proof>

lemma *summable-on-product-finite-right*:

fixes $f :: \langle 'a \times 'b \Rightarrow 'c :: \{ \text{topological-comm-monoid-add} \} \rangle$

assumes $\text{sum}: \langle \bigwedge y. y \in Y \Longrightarrow (\lambda x. f(x, y)) \text{ summable-on } X \rangle$

assumes $\langle \text{finite } Y \rangle$

shows $\langle f \text{ summable-on } (X \times Y) \rangle$

<proof>

2.6 Complex numbers

lemma *cmod-Re*:

assumes $x \geq 0$
shows $cmod\ x = Re\ x$
 $\langle proof \rangle$

lemma *abs-complex-real[simp]*: $abs\ x \in \mathbb{R}$ **for** $x :: complex$
 $\langle proof \rangle$

lemma *Im-abs[simp]*: $Im\ (abs\ x) = 0$
 $\langle proof \rangle$

lemma *cnj-x-x*: $cnj\ x * x = (abs\ x)^2$
 $\langle proof \rangle$

lemma *cnj-x-x-geq0[simp]*: $\langle cnj\ x * x \geq 0 \rangle$
 $\langle proof \rangle$

lemma *complex-of-real-leq-1-iff[iff]*: $\langle complex-of-real\ x \leq 1 \longleftrightarrow x \leq 1 \rangle$
 $\langle proof \rangle$

lemma *x-cnj-x*: $\langle x * cnj\ x = (abs\ x)^2 \rangle$
 $\langle proof \rangle$

lemma *has-sum-mono-neutral-complex*:
fixes $f :: 'a \Rightarrow complex$
assumes $\langle (f\ has-sum\ a)\ A \rangle$ **and** $\langle (g\ has-sum\ b)\ B \rangle$
assumes $\langle \bigwedge x. x \in A \cap B \implies f\ x \leq g\ x \rangle$
assumes $\langle \bigwedge x. x \in A - B \implies f\ x \leq 0 \rangle$
assumes $\langle \bigwedge x. x \in B - A \implies g\ x \geq 0 \rangle$
shows $a \leq b$
 $\langle proof \rangle$

lemma *Re-mono*: $x \leq y \implies Re\ x \leq Re\ y$
 $\langle proof \rangle$

lemma *comp-Im-same*: $x \leq y \implies Im\ x = Im\ y$
 $\langle proof \rangle$

lemma *Re-strict-mono*: $x < y \implies Re\ x < Re\ y$
 $\langle proof \rangle$

lemma *complex-of-real-cmod*: $\langle complex-of-real\ (cmod\ x) = abs\ x \rangle$
 $\langle proof \rangle$

lemma *less-eq-complexI*: $Re\ x \leq Re\ y \implies Im\ x = Im\ y \implies x \leq y$ $\langle proof \rangle$
lemma *less-complexI*: $Re\ x < Re\ y \implies Im\ x = Im\ y \implies x < y$ $\langle proof \rangle$

lemma *sums-le-complex*:
fixes $f\ g :: nat \Rightarrow complex$

```

assumes  $\langle \bigwedge n. f\ n \leq g\ n \rangle$ 
assumes  $\langle f\ \text{sums}\ s \rangle$ 
assumes  $\langle g\ \text{sums}\ t \rangle$ 
shows  $\langle s \leq t \rangle$ 
 $\langle \text{proof} \rangle$ 

```

2.7 List indices and enum

```

fun index-of where
  index-of  $x\ [] = (0::nat)$ 
| index-of  $x\ (y\#\text{ys}) = (\text{if } x=y \text{ then } 0 \text{ else } (\text{index-of } x\ \text{ys} + 1))$ 

```

definition *enum-idx* $(x::'a::\text{enum}) = \text{index-of } x\ (\text{enum-class.enum} :: 'a\ \text{list})$

```

lemma index-of-length: index-of  $x\ y \leq \text{length } y$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma index-of-correct:
  assumes  $x \in \text{set } y$ 
  shows  $y\ !\ \text{index-of } x\ y = x$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma enum-idx-correct:
  Enum.enum ! enum-idx  $i = i$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma index-of-bound:
  assumes  $y \neq []$  and  $x \in \text{set } y$ 
  shows index-of  $x\ y < \text{length } y$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma enum-idx-bound[simp]: enum-idx  $x < \text{CARD}('a)$  for  $x :: 'a::\text{enum}$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma index-of-nth:
  assumes distinct  $xs$ 
  assumes  $i < \text{length } xs$ 
  shows index-of  $(xs\ !\ i)\ xs = i$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma enum-idx-enum:
  assumes  $\langle i < \text{CARD}('a::\text{enum}) \rangle$ 
  shows  $\langle \text{enum-idx } (\text{enum-class.enum } !\ i :: 'a) = i \rangle$ 
 $\langle \text{proof} \rangle$ 

```

2.8 Filtering lists/sets

```

lemma map-filter-map: List.map-filter  $f\ (\text{map } g\ l) = \text{List.map-filter } (f\ o\ g)\ l$ 
 $\langle \text{proof} \rangle$ 

```

lemma *map-filter-Some* [simp]: $List.map-filter (\lambda x. Some (f x)) l = map f l$
 ⟨proof⟩

lemma *filter-Un*: $Set.filter f (x \cup y) = Set.filter f x \cup Set.filter f y$
 ⟨proof⟩

lemma *Set-filter-unchanged*: $Set.filter P X = X$ if $\bigwedge x. x \in X \implies P x$ for P and
 $X :: 'z set$
 ⟨proof⟩

2.9 Maps

definition *inj-map* $\pi = (\forall x y. \pi x = \pi y \wedge \pi x \neq None \longrightarrow x = y)$

definition *inv-map* $\pi = (\lambda y. \text{if } Some y \in range \pi \text{ then } Some (inv \pi (Some y)) \text{ else } None)$

lemma *inj-map-total*[simp]: $inj-map (Some \circ \pi) = inj \pi$
 ⟨proof⟩

lemma *inj-map-Some*[simp]: $inj-map Some$
 ⟨proof⟩

lemma *inv-map-total*:
 assumes *surj* π
 shows $inv-map (Some \circ \pi) = Some \circ inv \pi$
 ⟨proof⟩

lemma *inj-map-map-comp*[simp]:
 assumes *a1*: $inj-map f$ and *a2*: $inj-map g$
 shows $inj-map (f \circ_m g)$
 ⟨proof⟩

lemma *inj-map-inv-map*[simp]: $inj-map (inv-map \pi)$
 ⟨proof⟩

2.10 Lattices

unbundle *lattice-syntax*

The following lemma is identical to *Complete-Lattices.uminus-Inf* except for the more general sort.

lemma *uminus-Inf*: $- (\bigwedge A) = \bigsqcup (uminus \text{ ` } A)$ for $A :: \langle 'a :: complete-orthocomplemented-lattice set \rangle$
 ⟨proof⟩

The following lemma is identical to *Complete-Lattices.uminus-INF* except for the more general sort.

lemma *uminus-INF*: $\neg (\text{INF } x \in A. B \ x) = (\text{SUP } x \in A. \neg B \ x)$ **for** $B :: \langle 'a \Rightarrow 'b :: \text{complete-orthocomplemented-lattice} \rangle$
 $\langle \text{proof} \rangle$

The following lemma is identical to *Complete-Lattices.uminus-Sup* except for the more general sort.

lemma *uminus-Sup*: $\neg (\bigsqcup A) = \bigsqcap (\text{uminus } A)$ **for** $A :: \langle 'a :: \text{complete-orthocomplemented-lattice set} \rangle$
 $\langle \text{proof} \rangle$

The following lemma is identical to *Complete-Lattices.uminus-SUP* except for the more general sort.

lemma *uminus-SUP*: $\neg (\text{SUP } x \in A. B \ x) = (\text{INF } x \in A. \neg B \ x)$ **for** $B :: \langle 'a \Rightarrow 'b :: \text{complete-orthocomplemented-lattice} \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sumI-metric*:

fixes $l :: \langle 'a :: \{\text{metric-space, comm-monoid-add}\} \rangle$
assumes $\langle \bigwedge e. e > 0 \implies \exists X. \text{finite } X \wedge X \subseteq A \wedge (\forall Y. \text{finite } Y \wedge X \subseteq Y \wedge Y \subseteq A \longrightarrow \text{dist } (\text{sum } f \ Y) \ l < e) \rangle$
shows $\langle (f \text{ has-sum } l) \ A \rangle$
 $\langle \text{proof} \rangle$

lemma *limitin-pullback-topology*:

$\langle \text{limitin } (\text{pullback-topology } A \ g \ T) \ f \ l \ F \longleftrightarrow l \in A \wedge (\forall_F x \text{ in } F. f \ x \in A) \wedge \text{limitin } T \ (g \ o \ f) \ (g \ l) \ F \rangle$
 $\langle \text{proof} \rangle$

lemma *tendsto-coordinatewise*: $\langle (f \longrightarrow l) \ F \longleftrightarrow (\forall x. ((\lambda i. f \ i \ x) \longrightarrow l \ x) \ F) \rangle$
 $\langle \text{proof} \rangle$

lemma *limitin-closure-of*:

assumes *limit*: $\langle \text{limitin } T \ f \ c \ F \rangle$
assumes *in-S*: $\langle \forall_F x \text{ in } F. f \ x \in S \rangle$
assumes *nontrivial*: $\langle \neg \text{trivial-limit } F \rangle$
shows $\langle c \in T \text{ closure-of } S \rangle$
 $\langle \text{proof} \rangle$

2.11 Separating sets

definition *separating-set* :: $\langle (('a \Rightarrow 'b) \Rightarrow \text{bool}) \Rightarrow 'a \text{ set} \Rightarrow \text{bool} \rangle$ **where**
 $\langle \text{separating-set } P \ S \longleftrightarrow (\forall f \ g. P \ f \longrightarrow P \ g \longrightarrow (\forall x \in S. f \ x = g \ x) \longrightarrow f = g) \rangle$

lemma *separating-set-mono*: $\langle S \subseteq T \implies \text{separating-set } P \ S \implies \text{separating-set } P \ T \rangle$
 $\langle \text{proof} \rangle$

lemma *separating-set-UNIV[simp]*: $\langle \text{separating-set } P \ \text{UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *eq-from-separatingI*:
assumes $\langle \text{separating-set } P \ S \rangle$
assumes $\langle P \ f \rangle$ **and** $\langle P \ g \rangle$
assumes $\langle \bigwedge x. x \in S \implies f \ x = g \ x \rangle$
shows $\langle f = g \rangle$
 $\langle \text{proof} \rangle$

lemma *eq-from-separatingIx*:
— When using this as a rule, best instantiate x explicitly.
assumes $\langle \text{separating-set } P \ S \rangle$
assumes $\langle P \ f \rangle$ **and** $\langle P \ g \rangle$
assumes $\langle \bigwedge x. x \in S \implies f \ x = g \ x \rangle$
shows $\langle f \ x = g \ x \rangle$
 $\langle \text{proof} \rangle$

lemma *eq-from-separatingI2*:
assumes $\langle \text{separating-set } P \ ((\lambda(x,y). h \ x \ y) \ ^\circ (S \times T)) \rangle$
assumes $\langle P \ f \rangle$ **and** $\langle P \ g \rangle$
assumes $\langle \bigwedge x \ y. x \in S \implies y \in T \implies f \ (h \ x \ y) = g \ (h \ x \ y) \rangle$
shows $\langle f = g \rangle$
 $\langle \text{proof} \rangle$

lemma *eq-from-separatingI2x*:
— When using this as a rule, best instantiate x explicitly.
assumes $\langle \text{separating-set } P \ ((\lambda(x,y). h \ x \ y) \ ^\circ (S \times T)) \rangle$
assumes $\langle P \ f \rangle$ **and** $\langle P \ g \rangle$
assumes $\langle \bigwedge x \ y. x \in S \implies y \in T \implies f \ (h \ x \ y) = g \ (h \ x \ y) \rangle$
shows $\langle f \ x = g \ x \rangle$
 $\langle \text{proof} \rangle$

lemma *separating-setI*:
assumes $\langle \bigwedge f \ g. P \ f \implies P \ g \implies (\bigwedge x. x \in S \implies f \ x = g \ x) \implies f = g \rangle$
shows $\langle \text{separating-set } P \ S \rangle$
 $\langle \text{proof} \rangle$

end

3 *Extra-Vector-Spaces* – Additional facts about vector spaces

theory *Extra-Vector-Spaces*
imports
HOL-Analysis.Inner-Product
HOL-Analysis.Euclidean-Space
HOL-Library.Indicator-Function
HOL-Analysis.Topology-Euclidean-Space
HOL-Analysis.Line-Segment

begin

3.1 Euclidean spaces

typedef 'a euclidean-space = UNIV :: ('a $\Rightarrow$ real) set $\langle$ proof $\rangle$
setup-lifting type-definition-euclidean-space

instantiation euclidean-space :: (type) real-vector **begin**

lift-definition plus-euclidean-space ::

'a euclidean-space $\Rightarrow$ 'a euclidean-space $\Rightarrow$ 'a euclidean-space

is $\lambda f g x. f x + g x$ $\langle$ proof $\rangle$

lift-definition zero-euclidean-space :: 'a euclidean-space **is** $\lambda x. 0$ $\langle$ proof $\rangle$

lift-definition uminus-euclidean-space ::

'a euclidean-space $\Rightarrow$ 'a euclidean-space

is $\lambda f x. - f x$ $\langle$ proof $\rangle$

lift-definition minus-euclidean-space ::

'a euclidean-space $\Rightarrow$ 'a euclidean-space $\Rightarrow$ 'a euclidean-space

is $\lambda f g x. f x - g x$ $\langle$ proof $\rangle$

lift-definition scaleR-euclidean-space ::

real $\Rightarrow$ 'a euclidean-space $\Rightarrow$ 'a euclidean-space

is $\lambda c f x. c * f x$ $\langle$ proof $\rangle$

instance

$\langle$ proof $\rangle$

end

instantiation euclidean-space :: (finite) real-inner **begin**

lift-definition inner-euclidean-space :: 'a euclidean-space $\Rightarrow$ 'a euclidean-space $\Rightarrow$ real

is $\lambda f g. \sum x \in \text{UNIV}. f x * g x$:: real $\langle$ proof $\rangle$

definition norm-euclidean-space (x::'a euclidean-space) = sqrt (inner x x)

definition dist-euclidean-space (x::'a euclidean-space) y = norm (x-y)

definition sgn x = x /_R norm x **for** x::'a euclidean-space

definition uniformity = (INF e $\in$ {0<..}. principal {(x::'a euclidean-space, y). dist x y < e})

definition open U = ($\forall x \in U. \forall_F (x'::'a \text{ euclidean-space}, y) \text{ in uniformity. } x' = x \longrightarrow y \in U$)

instance

$\langle$ proof $\rangle$

end

instantiation euclidean-space :: (finite) euclidean-space **begin**

lift-definition euclidean-space-basis-vector :: 'a $\Rightarrow$ 'a euclidean-space **is**

$\lambda x. \text{indicator } \{x\}$ $\langle$ proof $\rangle$

definition Basis-euclidean-space == (euclidean-space-basis-vector ' UNIV)

instance

$\langle$ proof $\rangle$

end

3.2 Misc

lemma *closure-bounded-linear-image-subset-eq*:
assumes f : *bounded-linear* f
shows $\text{closure } (f \text{ ` closure } S) = \text{closure } (f \text{ ` } S)$
 $\langle \text{proof} \rangle$

lemma *not-singleton-real-normed-is-perfect-space[simp]*: $\langle \text{class.perfect-space } (\text{open} :: 'a :: \{\text{not-singleton}, \text{real-normed-vector}\} \text{ set} \Rightarrow \text{bool}) \rangle$
 $\langle \text{proof} \rangle$

lemma *infsum-bounded-linear*:
assumes $\langle \text{bounded-linear } h \rangle$
assumes $\langle f \text{ summable-on } A \rangle$
shows $\langle \text{infsum } (\lambda x. h (f x)) A = h (\text{infsum } f A) \rangle$
 $\langle \text{proof} \rangle$

lemma *abs-summable-on-bounded-linear*:
fixes $h f A$
assumes $\langle \text{bounded-linear } h \rangle$
assumes $\langle f \text{ abs-summable-on } A \rangle$
shows $\langle (h \circ f) \text{ abs-summable-on } A \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-plus-leq-norm-prod*: $\langle \text{norm } (a + b) \leq \text{sqrt } 2 * \text{norm } (a, b) \rangle$
 $\langle \text{proof} \rangle$

lemma *ex-norm1*:
assumes $\langle (UNIV :: 'a :: \text{real-normed-vector} \text{ set}) \neq \{0\} \rangle$
shows $\langle \exists x :: 'a. \text{norm } x = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *bdd-above-norm-f*:
assumes *bounded-linear* f
shows $\langle \text{bdd-above } \{\text{norm } (f x) \mid x. \text{norm } x = 1\} \rangle$
 $\langle \text{proof} \rangle$

lemma *any-norm-exists*:
assumes $\langle n \geq 0 \rangle$
shows $\langle \exists \psi :: 'a :: \{\text{real-normed-vector}, \text{not-singleton}\}. \text{norm } \psi = n \rangle$
 $\langle \text{proof} \rangle$

lemma *abs-summable-on-scaleR-left [intro]*:
fixes $c :: \langle 'a :: \text{real-normed-vector} \rangle$
assumes $c \neq 0 \implies f \text{ abs-summable-on } A$
shows $\langle \lambda x. f x *_R c \text{ abs-summable-on } A \rangle$
 $\langle \text{proof} \rangle$

lemma *abs-summable-on-scaleR-right [intro]*:

```

fixes  $f :: \langle 'a \Rightarrow 'b :: \text{real-normed-vector} \rangle$ 
assumes  $c \neq 0 \implies f \text{ abs-summable-on } A$ 
shows  $(\lambda x. c *_{\mathbb{R}} f x) \text{ abs-summable-on } A$ 
 $\langle \text{proof} \rangle$ 

```

lemma *ex-norm1-not-singleton*:

```

shows  $\langle \exists x :: 'a :: \{\text{real-normed-vector}, \text{not-singleton}\}. \text{norm } x = 1 \rangle$ 
 $\langle \text{proof} \rangle$ 

```

end

4 *Extra-Ordered-Fields* – Additional facts about ordered fields

theory *Extra-Ordered-Fields*

imports *Complex-Main HOL-Library.Complex-Order*
begin

4.1 Ordered Fields

In this section we introduce some type classes for ordered rings/fields/etc. that are weakenings of existing classes. Most theorems in this section are copies of the eponymous theorems from Isabelle/HOL, except that they are now proven requiring weaker type classes (usually the need for a total order is removed).

Since the lemmas are identical to the originals except for weaker type constraints, we use the same names as for the original lemmas. (In fact, the new lemmas could replace the original ones in Isabelle/HOL with at most minor incompatibilities.

4.2 Missing from Rings.thy

The existing class *abs-if* requires $|a| = (\text{if } a < 0 \text{ then } -a \text{ else } a)$. However, if $(<)$ is not a total order, this condition is too strong when a is incomparable with 0 . (Namely, it requires the absolute value to be the identity on such elements. E.g., the absolute value for complex numbers does not satisfy this.) The following class *partial-abs-if* is analogous to *abs-if* but does not require anything if a is incomparable with 0 .

```

class partial-abs-if = minus + uminus + ord + zero + abs +
  assumes abs-neg:  $a \leq 0 \implies \text{abs } a = -a$ 
  assumes abs-pos:  $a \geq 0 \implies \text{abs } a = a$ 

```

```

class ordered-ring-strict = ring + ordered-semiring-strict
  + ordered-ab-group-add + partial-abs-if
  — missing class analogous to linordered-ring-strict without requiring a total order

```

begin

subclass *ordered-ring* $\langle \text{proof} \rangle$

lemma *mult-strict-left-mono-neg*: $b < a \implies c < 0 \implies c * a < c * b$
 $\langle \text{proof} \rangle$

lemma *mult-strict-right-mono-neg*: $b < a \implies c < 0 \implies a * c < b * c$
 $\langle \text{proof} \rangle$

lemma *mult-neg-neg*: $a < 0 \implies b < 0 \implies 0 < a * b$
 $\langle \text{proof} \rangle$

end

lemmas *mult-sign-intros* =
 mult-nonneg-nonneg mult-nonneg-nonpos
 mult-nonpos-nonneg mult-nonpos-nonpos
 mult-pos-pos mult-pos-neg
 mult-neg-pos mult-neg-neg

4.3 Ordered fields

class *ordered-field* = *field* + *order* + *ordered-comm-semiring-strict* + *ordered-ab-group-add*
 + *partial-abs-if*
 — missing class analogous to *linordered-field* without requiring a total order

begin

lemma *frac-less-eq*:
 $y \neq 0 \implies z \neq 0 \implies x / y < w / z \longleftrightarrow (x * z - w * y) / (y * z) < 0$
 $\langle \text{proof} \rangle$

lemma *frac-le-eq*:
 $y \neq 0 \implies z \neq 0 \implies x / y \leq w / z \longleftrightarrow (x * z - w * y) / (y * z) \leq 0$
 $\langle \text{proof} \rangle$

lemmas *sign-simps* = *algebra-simps zero-less-mult-iff mult-less-0-iff*

lemmas (**in** $-$) *sign-simps* = *algebra-simps zero-less-mult-iff mult-less-0-iff*

Simplify expressions equated with 1

lemma *zero-eq-1-divide-iff* [*simp*]: $0 = 1 / a \longleftrightarrow a = 0$
 $\langle \text{proof} \rangle$

lemma *one-divide-eq-0-iff* [*simp*]: $1 / a = 0 \longleftrightarrow a = 0$
 $\langle \text{proof} \rangle$

Simplify expressions such as $0 < 1/x$ to $0 < x$

Simplify quotients that are compared with the value 1.

Conditional Simplification Rules: No Case Splits

lemma *eq-divide-eq-1* [*simp*]:
 $(1 = b/a) = ((a \neq 0 \ \& \ a = b))$
 $\langle \text{proof} \rangle$

lemma *divide-eq-eq-1* [*simp*]:
 $(b/a = 1) = ((a \neq 0 \ \& \ a = b))$
 $\langle \text{proof} \rangle$

end

The following type class intends to capture some important properties that are common both to the real and the complex numbers. The purpose is to be able to state and prove lemmas that apply both to the real and the complex numbers without needing to state the lemma twice.

class *nice-ordered-field* = *ordered-field* + *zero-less-one* + *idom-abs-sgn* +
assumes *positive-imp-inverse-positive*: $0 < a \implies 0 < \text{inverse } a$
and *inverse-le-imp-le*: $\text{inverse } a \leq \text{inverse } b \implies 0 < a \implies b \leq a$
and *dense-le*: $(\bigwedge x. x < y \implies x \leq z) \implies y \leq z$
and *nn-comparable*: $0 \leq a \implies 0 \leq b \implies a \leq b \vee b \leq a$
and *abs-nn*: $|x| \geq 0$

begin

subclass (in *linordered-field*) *nice-ordered-field*
 $\langle \text{proof} \rangle$

lemma *comparable*:
assumes *h1*: $a \leq c \vee a \geq c$
and *h2*: $b \leq c \vee b \geq c$
shows $a \leq b \vee b \leq a$
 $\langle \text{proof} \rangle$

lemma *negative-imp-inverse-negative*:
 $a < 0 \implies \text{inverse } a < 0$
 $\langle \text{proof} \rangle$

lemma *inverse-positive-imp-positive*:
assumes *inv-gt-0*: $0 < \text{inverse } a$ **and** *nz*: $a \neq 0$
shows $0 < a$
 $\langle \text{proof} \rangle$

lemma *inverse-negative-imp-negative*:
assumes *inv-less-0*: $\text{inverse } a < 0$ **and** *nz*: $a \neq 0$
shows $a < 0$
 $\langle \text{proof} \rangle$

lemma *linordered-field-no-lb*:
 $\forall x. \exists y. y < x$
 $\langle \text{proof} \rangle$

lemma *linordered-field-no-ub*:

$\forall x. \exists y. y > x$
 $\langle \text{proof} \rangle$

lemma *less-imp-inverse-less*:

assumes *less*: $a < b$ **and** *apos*: $0 < a$
shows $\text{inverse } b < \text{inverse } a$
 $\langle \text{proof} \rangle$

lemma *inverse-less-imp-less*:

$\text{inverse } a < \text{inverse } b \implies 0 < a \implies b < a$
 $\langle \text{proof} \rangle$

Both premises are essential. Consider -1 and 1.

lemma *inverse-less-iff-less* [simp]:

$0 < a \implies 0 < b \implies \text{inverse } a < \text{inverse } b \longleftrightarrow b < a$
 $\langle \text{proof} \rangle$

lemma *le-imp-inverse-le*:

$a \leq b \implies 0 < a \implies \text{inverse } b \leq \text{inverse } a$
 $\langle \text{proof} \rangle$

lemma *inverse-le-iff-le* [simp]:

$0 < a \implies 0 < b \implies \text{inverse } a \leq \text{inverse } b \longleftrightarrow b \leq a$
 $\langle \text{proof} \rangle$

These results refer to both operands being negative. The opposite-sign case is trivial, since inverse preserves signs.

lemma *inverse-le-imp-le-neg*:

$\text{inverse } a \leq \text{inverse } b \implies b < 0 \implies b \leq a$
 $\langle \text{proof} \rangle$

lemma *inverse-less-imp-less-neg*:

$\text{inverse } a < \text{inverse } b \implies b < 0 \implies b < a$
 $\langle \text{proof} \rangle$

lemma *inverse-less-iff-less-neg* [simp]:

$a < 0 \implies b < 0 \implies \text{inverse } a < \text{inverse } b \longleftrightarrow b < a$
 $\langle \text{proof} \rangle$

lemma *le-imp-inverse-le-neg*:

$a \leq b \implies b < 0 \implies \text{inverse } b \leq \text{inverse } a$
 $\langle \text{proof} \rangle$

lemma *inverse-le-iff-le-neg* [simp]:

$a < 0 \implies b < 0 \implies \text{inverse } a \leq \text{inverse } b \longleftrightarrow b \leq a$
 $\langle \text{proof} \rangle$

lemma *one-less-inverse*:
 $0 < a \implies a < 1 \implies 1 < \text{inverse } a$
 $\langle \text{proof} \rangle$

lemma *one-le-inverse*:
 $0 < a \implies a \leq 1 \implies 1 \leq \text{inverse } a$
 $\langle \text{proof} \rangle$

lemma *pos-le-divide-eq* [*field-simps*]:
assumes $0 < c$
shows $a \leq b / c \longleftrightarrow a * c \leq b$
 $\langle \text{proof} \rangle$

lemma *pos-less-divide-eq* [*field-simps*]:
assumes $0 < c$
shows $a < b / c \longleftrightarrow a * c < b$
 $\langle \text{proof} \rangle$

lemma *neg-less-divide-eq* [*field-simps*]:
assumes $c < 0$
shows $a < b / c \longleftrightarrow b < a * c$
 $\langle \text{proof} \rangle$

lemma *neg-le-divide-eq* [*field-simps*]:
assumes $c < 0$
shows $a \leq b / c \longleftrightarrow b \leq a * c$
 $\langle \text{proof} \rangle$

lemma *pos-divide-le-eq* [*field-simps*]:
assumes $0 < c$
shows $b / c \leq a \longleftrightarrow b \leq a * c$
 $\langle \text{proof} \rangle$

lemma *pos-divide-less-eq* [*field-simps*]:
assumes $0 < c$
shows $b / c < a \longleftrightarrow b < a * c$
 $\langle \text{proof} \rangle$

lemma *neg-divide-le-eq* [*field-simps*]:
assumes $c < 0$
shows $b / c \leq a \longleftrightarrow a * c \leq b$
 $\langle \text{proof} \rangle$

lemma *neg-divide-less-eq* [*field-simps*]:
assumes $c < 0$
shows $b / c < a \longleftrightarrow a * c < b$
 $\langle \text{proof} \rangle$

The following *field-simps* rules are necessary, as minus is always moved atop

of division but we want to get rid of division.

lemma *pos-le-minus-divide-eq* [*field-simps*]: $0 < c \implies a \leq - (b / c) \longleftrightarrow a * c \leq - b$
 ⟨*proof*⟩

lemma *neg-le-minus-divide-eq* [*field-simps*]: $c < 0 \implies a \leq - (b / c) \longleftrightarrow - b \leq a * c$
 ⟨*proof*⟩

lemma *pos-less-minus-divide-eq* [*field-simps*]: $0 < c \implies a < - (b / c) \longleftrightarrow a * c < - b$
 ⟨*proof*⟩

lemma *neg-less-minus-divide-eq* [*field-simps*]: $c < 0 \implies a < - (b / c) \longleftrightarrow - b < a * c$
 ⟨*proof*⟩

lemma *pos-minus-divide-less-eq* [*field-simps*]: $0 < c \implies - (b / c) < a \longleftrightarrow - b < a * c$
 ⟨*proof*⟩

lemma *neg-minus-divide-less-eq* [*field-simps*]: $c < 0 \implies - (b / c) < a \longleftrightarrow a * c < - b$
 ⟨*proof*⟩

lemma *pos-minus-divide-le-eq* [*field-simps*]: $0 < c \implies - (b / c) \leq a \longleftrightarrow - b \leq a * c$
 ⟨*proof*⟩

lemma *neg-minus-divide-le-eq* [*field-simps*]: $c < 0 \implies - (b / c) \leq a \longleftrightarrow a * c \leq - b$
 ⟨*proof*⟩

lemma *frac-less-eq*:
 $y \neq 0 \implies z \neq 0 \implies x / y < w / z \longleftrightarrow (x * z - w * y) / (y * z) < 0$
 ⟨*proof*⟩

lemma *frac-le-eq*:
 $y \neq 0 \implies z \neq 0 \implies x / y \leq w / z \longleftrightarrow (x * z - w * y) / (y * z) \leq 0$
 ⟨*proof*⟩

Lemmas *sign-simps* is a first attempt to automate proofs of positivity/negativity needed for *field-simps*. Have not added *sign-simps* to *field-simps* because the former can lead to case explosions.

lemma *divide-pos-pos[simp]*:
 $0 < x \implies 0 < y \implies 0 < x / y$
 ⟨*proof*⟩

lemma *divide-nonneg-pos*:

$$0 \leq x \implies 0 < y \implies 0 \leq x / y$$

<proof>

lemma *divide-neg-pos:*

$$x < 0 \implies 0 < y \implies x / y < 0$$

<proof>

lemma *divide-nonpos-pos:*

$$x \leq 0 \implies 0 < y \implies x / y \leq 0$$

<proof>

lemma *divide-pos-neg:*

$$0 < x \implies y < 0 \implies x / y < 0$$

<proof>

lemma *divide-nonneg-neg:*

$$0 \leq x \implies y < 0 \implies x / y \leq 0$$

<proof>

lemma *divide-neg-neg:*

$$x < 0 \implies y < 0 \implies 0 < x / y$$

<proof>

lemma *divide-nonpos-neg:*

$$x \leq 0 \implies y < 0 \implies 0 \leq x / y$$

<proof>

lemma *divide-strict-right-mono:*

$$a < b \implies 0 < c \implies a / c < b / c$$

<proof>

lemma *divide-strict-right-mono-neg:*

$$b < a \implies c < 0 \implies a / c < b / c$$

<proof>

The last premise ensures that a and b have the same sign

lemma *divide-strict-left-mono:*

$$b < a \implies 0 < c \implies 0 < a*b \implies c / a < c / b$$

<proof>

lemma *divide-left-mono:*

$$b \leq a \implies 0 \leq c \implies 0 < a*b \implies c / a \leq c / b$$

<proof>

lemma *divide-strict-left-mono-neg:*

$$a < b \implies c < 0 \implies 0 < a*b \implies c / a < c / b$$

<proof>

lemma *mult-imp-div-pos-le*: $0 < y \implies x \leq z * y \implies x / y \leq z$
 ⟨proof⟩

lemma *mult-imp-le-div-pos*: $0 < y \implies z * y \leq x \implies z \leq x / y$
 ⟨proof⟩

lemma *mult-imp-div-pos-less*: $0 < y \implies x < z * y \implies x / y < z$
 ⟨proof⟩

lemma *mult-imp-less-div-pos*: $0 < y \implies z * y < x \implies z < x / y$
 ⟨proof⟩

lemma *frac-le*: $0 \leq x \implies x \leq y \implies 0 < w \implies w \leq z \implies x / z \leq y / w$
 ⟨proof⟩

lemma *frac-less*: $0 \leq x \implies x < y \implies 0 < w \implies w \leq z \implies x / z < y / w$
 ⟨proof⟩

lemma *frac-less2*: $0 < x \implies x \leq y \implies 0 < w \implies w < z \implies x / z < y / w$
 ⟨proof⟩

lemma *less-half-sum*: $a < b \implies a < (a+b) / (1+1)$
 ⟨proof⟩

lemma *gt-half-sum*: $a < b \implies (a+b)/(1+1) < b$
 ⟨proof⟩

subclass *unbounded-dense-order*
 ⟨proof⟩

lemma *dense-le-bounded*:
 fixes $x\ y\ z :: 'a$
 assumes $x < y$
 and *: $\bigwedge w. \llbracket x < w ; w < y \rrbracket \implies w \leq z$
 shows $y \leq z$
 ⟨proof⟩

subclass *field-abs-sgn* ⟨proof⟩

lemma *nonzero-abs-inverse*:
 $a \neq 0 \implies |\text{inverse } a| = \text{inverse } |a|$
 ⟨proof⟩

lemma *nonzero-abs-divide*:
 $b \neq 0 \implies |a / b| = |a| / |b|$
 ⟨proof⟩

lemma *field-le-epsilon*:

assumes e : $\bigwedge e. 0 < e \implies x \leq y + e$

shows $x \leq y$

$\langle proof \rangle$

lemma *inverse-positive-iff-positive* [simp]:

$(0 < \text{inverse } a) = (0 < a)$

$\langle proof \rangle$

lemma *inverse-negative-iff-negative* [simp]:

$(\text{inverse } a < 0) = (a < 0)$

$\langle proof \rangle$

lemma *inverse-nonnegative-iff-nonnegative* [simp]:

$0 \leq \text{inverse } a \longleftrightarrow 0 \leq a$

$\langle proof \rangle$

lemma *inverse-nonpositive-iff-nonpositive* [simp]:

$\text{inverse } a \leq 0 \longleftrightarrow a \leq 0$

$\langle proof \rangle$

lemma *one-less-inverse-iff*: $1 < \text{inverse } x \longleftrightarrow 0 < x \wedge x < 1$

$\langle proof \rangle$

lemma *one-le-inverse-iff*: $1 \leq \text{inverse } x \longleftrightarrow 0 < x \wedge x \leq 1$

$\langle proof \rangle$

lemma *inverse-less-1-iff*: $\text{inverse } x < 1 \longleftrightarrow x \leq 0 \vee 1 < x$

$\langle proof \rangle$

lemma *inverse-le-1-iff*: $\text{inverse } x \leq 1 \longleftrightarrow x \leq 0 \vee 1 \leq x$

$\langle proof \rangle$

Simplify expressions such as $0 < 1/x$ to $0 < x$

lemma *zero-le-divide-1-iff* [simp]:

$0 \leq 1 / a \longleftrightarrow 0 \leq a$

$\langle proof \rangle$

lemma *zero-less-divide-1-iff* [simp]:

$0 < 1 / a \longleftrightarrow 0 < a$

$\langle proof \rangle$

lemma *divide-le-0-1-iff* [simp]:

$1 / a \leq 0 \longleftrightarrow a \leq 0$

$\langle proof \rangle$

lemma *divide-less-0-1-iff* [simp]:

$1 / a < 0 \longleftrightarrow a < 0$

$\langle proof \rangle$

lemma *divide-right-mono*:

$$a \leq b \implies 0 \leq c \implies a/c \leq b/c$$

$\langle proof \rangle$

lemma *divide-right-mono-neg*: $a \leq b$

$$\implies c \leq 0 \implies b / c \leq a / c$$

$\langle proof \rangle$

lemma *divide-left-mono-neg*: $a \leq b$

$$\implies c \leq 0 \implies 0 < a * b \implies c / a \leq c / b$$

$\langle proof \rangle$

lemma *divide-nonneg-nonneg* [simp]:

$$0 \leq x \implies 0 \leq y \implies 0 \leq x / y$$

$\langle proof \rangle$

lemma *divide-nonpos-nonpos*:

$$x \leq 0 \implies y \leq 0 \implies 0 \leq x / y$$

$\langle proof \rangle$

lemma *divide-nonneg-nonpos*:

$$0 \leq x \implies y \leq 0 \implies x / y \leq 0$$

$\langle proof \rangle$

lemma *divide-nonpos-nonneg*:

$$x \leq 0 \implies 0 \leq y \implies x / y \leq 0$$

$\langle proof \rangle$

Conditional Simplification Rules: No Case Splits

lemma *le-divide-eq-1-pos* [simp]:

$$0 < a \implies (1 \leq b/a) = (a \leq b)$$

$\langle proof \rangle$

lemma *le-divide-eq-1-neg* [simp]:

$$a < 0 \implies (1 \leq b/a) = (b \leq a)$$

$\langle proof \rangle$

lemma *divide-le-eq-1-pos* [simp]:

$$0 < a \implies (b/a \leq 1) = (b \leq a)$$

$\langle proof \rangle$

lemma *divide-le-eq-1-neg* [simp]:

$$a < 0 \implies (b/a \leq 1) = (a \leq b)$$

$\langle proof \rangle$

lemma *less-divide-eq-1-pos* [simp]:

$$0 < a \implies (1 < b/a) = (a < b)$$

$\langle proof \rangle$

lemma *less-divide-eq-1-neg* [simp]:
 $a < 0 \implies (1 < b/a) = (b < a)$
 $\langle proof \rangle$

lemma *divide-less-eq-1-pos* [simp]:
 $0 < a \implies (b/a < 1) = (b < a)$
 $\langle proof \rangle$

lemma *divide-less-eq-1-neg* [simp]:
 $a < 0 \implies b/a < 1 \longleftrightarrow a < b$
 $\langle proof \rangle$

lemma *abs-div-pos*: $0 < y \implies$
 $|x| / y = |x / y|$
 $\langle proof \rangle$

lemma *zero-le-divide-abs-iff* [simp]: $(0 \leq a / |b|) = (0 \leq a \mid b = 0)$
 $\langle proof \rangle$

lemma *divide-le-0-abs-iff* [simp]: $(a / |b| \leq 0) = (a \leq 0 \mid b = 0)$
 $\langle proof \rangle$

For creating values between u and v .

lemma *scaling-mono*:
assumes $u \leq v$ **and** $0 \leq r$ **and** $r \leq s$
shows $u + r * (v - u) / s \leq v$
 $\langle proof \rangle$

end

code-identifier

code-module *Ordered-Fields* $\mapsto$ (*SML*) *Arith* **and** (*OCaml*) *Arith* **and** (*Haskell*)
Arith

4.4 Ordering on complex numbers

instantiation *complex* :: *nice-ordered-field* **begin**
instance
 $\langle proof \rangle$
end

lemma *complex-of-real-mono*:
 $x \leq y \implies \text{complex-of-real } x \leq \text{complex-of-real } y$
 $\langle proof \rangle$

lemma *complex-of-real-mono-iff*[simp]:
 $\text{complex-of-real } x \leq \text{complex-of-real } y \longleftrightarrow x \leq y$
 ⟨proof⟩

lemma *complex-of-real-strict-mono-iff*[simp]:
 $\text{complex-of-real } x < \text{complex-of-real } y \longleftrightarrow x < y$
 ⟨proof⟩

lemma *complex-of-real-nn-iff*[simp]:
 $0 \leq \text{complex-of-real } y \longleftrightarrow 0 \leq y$
 ⟨proof⟩

lemma *complex-of-real-pos-iff*[simp]:
 $0 < \text{complex-of-real } y \longleftrightarrow 0 < y$
 ⟨proof⟩

end

5 *Extra-Operator-Norm* – Additional facts bout the operator norm

theory *Extra-Operator-Norm*
imports *HOL–Analysis.Operator-Norm*
Extra-General
HOL–Analysis.Bounded-Linear-Function
Extra-Vector-Spaces
begin

This theorem complements *HOL–Analysis.Operator-Norm* additional useful facts about operator norms.

lemma *onorm-sphere*:
fixes $f :: 'a::\{\text{real-normed-vector}, \text{not-singleton}\} \Rightarrow 'b::\text{real-normed-vector}$
assumes $a1: \text{bounded-linear } f$
shows $\langle \text{onorm } f = \text{Sup } \{\text{norm } (f x) \mid x. \text{norm } x = 1\} \rangle$
 ⟨proof⟩

lemma *onormI*:
assumes $\bigwedge x. \text{norm } (f x) \leq b * \text{norm } x$
and $x \neq 0$ **and** $\text{norm } (f x) = b * \text{norm } x$
shows $\text{onorm } f = b$
 ⟨proof⟩

end

6 *Complex-Vector-Spaces0* – Vector Spaces and Algebras over the Complex Numbers

```

theory Complex-Vector-Spaces0
  imports HOL.Real-Vector-Spaces HOL.Topological-Spaces HOL.Vector-Spaces
    Complex-Main
    HOL-Library.Complex-Order
    HOL-Analysis.Product-Vector
begin

```

6.1 Complex vector spaces

```

class scaleC = scaleR +
  fixes scaleC :: complex  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr ' *_C ' 75)
  assumes scaleR-scaleC: scaleR r = scaleC (complex-of-real r)
begin

abbreviation divideC :: 'a  $\Rightarrow$  complex  $\Rightarrow$  'a (infixl ' /_C ' 70)
  where x /_C c  $\equiv$  inverse c *_C x

end

class complex-vector = scaleC + ab-group-add +
  assumes scaleC-add-right: a *_C (x + y) = (a *_C x) + (a *_C y)
  and scaleC-add-left: (a + b) *_C x = (a *_C x) + (b *_C x)
  and scaleC-scaleC[simp]: a *_C (b *_C x) = (a * b) *_C x
  and scaleC-one[simp]: 1 *_C x = x

subclass (in complex-vector) real-vector
  <proof>

class complex-algebra = complex-vector + ring +
  assumes mult-scaleC-left[simp]: a *_C x * y = a *_C (x * y)
  and mult-scaleC-right[simp]: x * a *_C y = a *_C (x * y)

subclass (in complex-algebra) real-algebra
  <proof>

class complex-algebra-1 = complex-algebra + ring-1

subclass (in complex-algebra-1) real-algebra-1 <proof>

class complex-div-algebra = complex-algebra-1 + division-ring

subclass (in complex-div-algebra) real-div-algebra <proof>

```

```

class complex-field = complex-div-algebra + field

subclass (in complex-field) real-field ⟨proof⟩

instantiation complex :: complex-field
begin

definition complex-scaleC-def [simp]: scaleC a x = a * x

instance
  ⟨proof⟩

end

locale clinear = Vector-Spaces.linear scaleC::-=>-=>'a::complex-vector scaleC::-=>-=>'b::complex-vector
begin

sublocale real: linear
  — Gives access to all lemmas from Real-Vector-Spaces.linear using prefix real.
  ⟨proof⟩

lemmas scaleC = scale

end

global-interpretation complex-vector: vector-space scaleC :: complex => 'a => 'a
:: complex-vector
  rewrites Vector-Spaces.linear (*C) (*C) = clinear
  and Vector-Spaces.linear (*) (*C) = clinear
  defines cdependent-raw-def: cdependent = complex-vector.dependent
  and crepresentation-raw-def: crepresentation = complex-vector.representation
  and csubspace-raw-def: csubspace = complex-vector.subspace
  and cspan-raw-def: cspan = complex-vector.span
  and cextend-basis-raw-def: cextend-basis = complex-vector.extend-basis
  and cdim-raw-def: cdim = complex-vector.dim
  ⟨proof⟩

abbreviation cindependent x ≡ ¬ cdependent x

global-interpretation complex-vector: vector-space-pair scaleC::-=>-=>'a::complex-vector
scaleC::-=>-=>'b::complex-vector
  rewrites Vector-Spaces.linear (*C) (*C) = clinear
  and Vector-Spaces.linear (*) (*C) = clinear

```


defines *cconstruct-raw-def*: *cconstruct* = *complex-vector.construct*
 <proof>

lemma *clinear-compose*: *clinear f* $\implies$ *clinear g* $\implies$ *clinear (g o f)*
 <proof>

Recover original theorem names

lemmas *scaleC-left-commute* = *complex-vector.scale-left-commute*
lemmas *scaleC-zero-left* = *complex-vector.scale-zero-left*
lemmas *scaleC-minus-left* = *complex-vector.scale-minus-left*
lemmas *scaleC-diff-left* = *complex-vector.scale-left-diff-distrib*
lemmas *scaleC-sum-left* = *complex-vector.scale-sum-left*
lemmas *scaleC-zero-right* = *complex-vector.scale-zero-right*
lemmas *scaleC-minus-right* = *complex-vector.scale-minus-right*
lemmas *scaleC-diff-right* = *complex-vector.scale-right-diff-distrib*
lemmas *scaleC-sum-right* = *complex-vector.scale-sum-right*
lemmas *scaleC-eq-0-iff* = *complex-vector.scale-eq-0-iff*
lemmas *scaleC-left-imp-eq* = *complex-vector.scale-left-imp-eq*
lemmas *scaleC-right-imp-eq* = *complex-vector.scale-right-imp-eq*
lemmas *scaleC-cancel-left* = *complex-vector.scale-cancel-left*
lemmas *scaleC-cancel-right* = *complex-vector.scale-cancel-right*

lemma *divideC-field-simps[field-simps]*:
 $c \neq 0 \implies a = b /_C c \iff c *_C a = b$
 $c \neq 0 \implies b /_C c = a \iff b = c *_C a$
 $c \neq 0 \implies a + b /_C c = (c *_C a + b) /_C c$
 $c \neq 0 \implies a /_C c + b = (a + c *_C b) /_C c$
 $c \neq 0 \implies a - b /_C c = (c *_C a - b) /_C c$
 $c \neq 0 \implies a /_C c - b = (a - c *_C b) /_C c$
 $c \neq 0 \implies -(a /_C c) + b = (-a + c *_C b) /_C c$
 $c \neq 0 \implies -(a /_C c) - b = (-a - c *_C b) /_C c$
for *a b* :: '*a*' :: *complex-vector*
 <proof>

Legacy names – omitted

lemmas *clinear-injective-0* = *linear-inj-iff-eq-0*
and *clinear-injective-on-subspace-0* = *linear-inj-on-iff-eq-0*
and *clinear-cmul* = *linear-scale*
and *clinear-scaleC* = *linear-scale-self*
and *csubspace-mul* = *subspace-scale*
and *cspan-linear-image* = *linear-span-image*
and *cspan-0* = *span-zero*
and *cspan-mul* = *span-scale*
and *injective-scaleC* = *injective-scale*

lemma *scaleC-minus1-left [simp]*: *scaleC* (-1) *x* = $- x$
for *x* :: '*a*' :: *complex-vector*

$\langle proof \rangle$

lemma *scaleC-2*:

fixes $x :: 'a::complex-vector$

shows $scaleC\ 2\ x = x + x$

$\langle proof \rangle$

lemma *scaleC-half-double* [*simp*]:

fixes $a :: 'a::complex-vector$

shows $(1 / 2) *_{\mathbb{C}} (a + a) = a$

$\langle proof \rangle$

lemma *clinear-scale-complex*:

fixes $c::complex$ **shows** $clinear\ f \implies f\ (c * b) = c * f\ b$

$\langle proof \rangle$

interpretation *scaleC-left*: *additive* $(\lambda a. scaleC\ a\ x :: 'a::complex-vector)$

$\langle proof \rangle$

interpretation *scaleC-right*: *additive* $(\lambda x. scaleC\ a\ x :: 'a::complex-vector)$

$\langle proof \rangle$

lemma *nonzero-inverse-scaleC-distrib*:

$a \neq 0 \implies x \neq 0 \implies inverse\ (scaleC\ a\ x) = scaleC\ (inverse\ a)\ (inverse\ x)$

for $x :: 'a::complex-div-algebra$

$\langle proof \rangle$

lemma *inverse-scaleC-distrib*: $inverse\ (scaleC\ a\ x) = scaleC\ (inverse\ a)\ (inverse\ x)$

for $x :: 'a::\{complex-div-algebra, division-ring\}$

$\langle proof \rangle$

lemma *complex-add-divide-simps*[*vector-add-divide-simps*]:

$v + (b / z) *_{\mathbb{C}} w = (if\ z = 0\ then\ v\ else\ (z *_{\mathbb{C}} v + b *_{\mathbb{C}} w) /_{\mathbb{C}} z)$

$a *_{\mathbb{C}} v + (b / z) *_{\mathbb{C}} w = (if\ z = 0\ then\ a *_{\mathbb{C}} v\ else\ ((a * z) *_{\mathbb{C}} v + b *_{\mathbb{C}} w) /_{\mathbb{C}} z)$

$(a / z) *_{\mathbb{C}} v + w = (if\ z = 0\ then\ w\ else\ (a *_{\mathbb{C}} v + z *_{\mathbb{C}} w) /_{\mathbb{C}} z)$

$(a / z) *_{\mathbb{C}} v + b *_{\mathbb{C}} w = (if\ z = 0\ then\ b *_{\mathbb{C}} w\ else\ (a *_{\mathbb{C}} v + (b * z) *_{\mathbb{C}} w) /_{\mathbb{C}} z)$

$v - (b / z) *_{\mathbb{C}} w = (if\ z = 0\ then\ v\ else\ (z *_{\mathbb{C}} v - b *_{\mathbb{C}} w) /_{\mathbb{C}} z)$

$a *_{\mathbb{C}} v - (b / z) *_{\mathbb{C}} w = (if\ z = 0\ then\ a *_{\mathbb{C}} v\ else\ ((a * z) *_{\mathbb{C}} v - b *_{\mathbb{C}} w) /_{\mathbb{C}} z)$

$(a / z) *_{\mathbb{C}} v - w = (if\ z = 0\ then\ -w\ else\ (a *_{\mathbb{C}} v - z *_{\mathbb{C}} w) /_{\mathbb{C}} z)$

$(a / z) *_{\mathbb{C}} v - b *_{\mathbb{C}} w = (if\ z = 0\ then\ -b *_{\mathbb{C}} w\ else\ (a *_{\mathbb{C}} v - (b * z) *_{\mathbb{C}} w) /_{\mathbb{C}} z)$

$/_C z)$
for $v :: 'a :: \text{complex-vector}$
 $\langle \text{proof} \rangle$

lemma *ceq-vector-fraction-iff* [vector-add-divide-simps]:
fixes $x :: 'a :: \text{complex-vector}$
shows $(x = (u / v) *_C a) \longleftrightarrow (\text{if } v=0 \text{ then } x = 0 \text{ else } v *_C x = u *_C a)$
 $\langle \text{proof} \rangle$

lemma *cvector-fraction-eq-iff* [vector-add-divide-simps]:
fixes $x :: 'a :: \text{complex-vector}$
shows $((u / v) *_C a = x) \longleftrightarrow (\text{if } v=0 \text{ then } x = 0 \text{ else } u *_C a = v *_C x)$
 $\langle \text{proof} \rangle$

lemma *complex-vector-affinity-eq*:
fixes $x :: 'a :: \text{complex-vector}$
assumes $m0: m \neq 0$
shows $m *_C x + c = y \longleftrightarrow x = \text{inverse } m *_C y - (\text{inverse } m *_C c)$
(is ?lhs $\longleftrightarrow$?rhs)
 $\langle \text{proof} \rangle$

lemma *complex-vector-eq-affinity*: $m \neq 0 \implies y = m *_C x + c \longleftrightarrow \text{inverse } m *_C y - (\text{inverse } m *_C c) = x$
for $x :: 'a :: \text{complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-eq-iff* [simp]: $b + u *_C a = a + u *_C b \longleftrightarrow a = b \vee u = 1$
for $a :: 'a :: \text{complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-collapse* [simp]: $(1 - u) *_C a + u *_C a = a$
for $a :: 'a :: \text{complex-vector}$
 $\langle \text{proof} \rangle$

6.2 Embedding of the Complex Numbers into any *complex-algebra-1*: *of-complex*

definition *of-complex* :: $\text{complex} \Rightarrow 'a :: \text{complex-algebra-1}$
where *of-complex* $c = \text{scaleC } c \ 1$

lemma *scaleC-conv-of-complex*: $\text{scaleC } r \ x = \text{of-complex } r * x$
 $\langle \text{proof} \rangle$

lemma *of-complex-0* [simp]: $\text{of-complex } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *of-complex-1* [simp]: $\text{of-complex } 1 = 1$
 $\langle \text{proof} \rangle$

lemma *of-complex-add* [simp]: *of-complex* $(x + y) = \textit{of-complex } x + \textit{of-complex } y$
 ⟨proof⟩

lemma *of-complex-minus* [simp]: *of-complex* $(- x) = - \textit{of-complex } x$
 ⟨proof⟩

lemma *of-complex-diff* [simp]: *of-complex* $(x - y) = \textit{of-complex } x - \textit{of-complex } y$
 ⟨proof⟩

lemma *of-complex-mult* [simp]: *of-complex* $(x * y) = \textit{of-complex } x * \textit{of-complex } y$
 ⟨proof⟩

lemma *of-complex-sum*[simp]: *of-complex* $(\textit{sum } f s) = (\sum_{x \in s}. \textit{of-complex } (f x))$
 ⟨proof⟩

lemma *of-complex-prod*[simp]: *of-complex* $(\textit{prod } f s) = (\prod_{x \in s}. \textit{of-complex } (f x))$
 ⟨proof⟩

lemma *nonzero-of-complex-inverse*:
 $x \neq 0 \implies \textit{of-complex } (\textit{inverse } x) = \textit{inverse } (\textit{of-complex } x :: 'a::\textit{complex-div-algebra})$
 ⟨proof⟩

lemma *of-complex-inverse* [simp]:
 $\textit{of-complex } (\textit{inverse } x) = \textit{inverse } (\textit{of-complex } x :: 'a::\{\textit{complex-div-algebra}, \textit{division-ring}\})$
 ⟨proof⟩

lemma *nonzero-of-complex-divide*:
 $y \neq 0 \implies \textit{of-complex } (x / y) = (\textit{of-complex } x / \textit{of-complex } y :: 'a::\textit{complex-field})$
 ⟨proof⟩

lemma *of-complex-divide* [simp]:
 $\textit{of-complex } (x / y) = (\textit{of-complex } x / \textit{of-complex } y :: 'a::\textit{complex-div-algebra})$
 ⟨proof⟩

lemma *of-complex-power* [simp]:
 $\textit{of-complex } (x ^ n) = (\textit{of-complex } x :: 'a::\{\textit{complex-algebra-1}\}) ^ n$
 ⟨proof⟩

lemma *of-complex-power-int* [simp]:
 $\textit{of-complex } (\textit{power-int } x n) = \textit{power-int } (\textit{of-complex } x :: 'a :: \{\textit{complex-div-algebra}, \textit{division-ring}\})$
 n
 ⟨proof⟩

lemma *of-complex-eq-iff* [simp]: *of-complex* $x = \textit{of-complex } y \longleftrightarrow x = y$
 ⟨proof⟩

lemma *inj-of-complex*: *inj of-complex*
 ⟨proof⟩

lemmas *of-complex-eq-0-iff* [simp] = *of-complex-eq-iff* [of - 0, simplified]

lemmas *of-complex-eq-1-iff* [simp] = *of-complex-eq-iff* [of - 1, simplified]

lemma *minus-of-complex-eq-of-complex-iff* [simp]: $-of-complex\ x = of-complex\ y$
 $\longleftrightarrow -x = y$
 ⟨proof⟩

lemma *of-complex-eq-minus-of-complex-iff* [simp]: $of-complex\ x = -of-complex\ y$
 $\longleftrightarrow x = -y$
 ⟨proof⟩

lemma *of-complex-eq-id* [simp]: $of-complex = (id :: complex \Rightarrow complex)$
 ⟨proof⟩

Collapse nested embeddings.

lemma *of-complex-of-nat-eq* [simp]: $of-complex\ (of-nat\ n) = of-nat\ n$
 ⟨proof⟩

lemma *of-complex-of-int-eq* [simp]: $of-complex\ (of-int\ z) = of-int\ z$
 ⟨proof⟩

lemma *of-complex-numeral* [simp]: $of-complex\ (numeral\ w) = numeral\ w$
 ⟨proof⟩

lemma *of-complex-neg-numeral* [simp]: $of-complex\ (-\ numeral\ w) = -\ numeral\ w$
 ⟨proof⟩

lemma *numeral-power-int-eq-of-complex-cancel-iff* [simp]:
 $power-int\ (numeral\ x)\ n = (of-complex\ y :: 'a :: \{complex-div-algebra, division-ring\}) \longleftrightarrow$
 $power-int\ (numeral\ x)\ n = y$
 ⟨proof⟩

lemma *of-complex-eq-numeral-power-int-cancel-iff* [simp]:
 $(of-complex\ y :: 'a :: \{complex-div-algebra, division-ring\}) = power-int\ (numeral\ x)\ n \longleftrightarrow$
 $y = power-int\ (numeral\ x)\ n$
 ⟨proof⟩

lemma *of-complex-eq-of-complex-power-int-cancel-iff* [simp]:
 $power-int\ (of-complex\ b :: 'a :: \{complex-div-algebra, division-ring\})\ w = of-complex\ x \longleftrightarrow$
 $power-int\ b\ w = x$
 ⟨proof⟩

lemma *of-complex-in-Ints-iff* [simp]: $of-complex\ x \in \mathbb{Z} \longleftrightarrow x \in \mathbb{Z}$
 ⟨proof⟩

lemma *Ints-of-complex* [intro]: $x \in \mathbb{Z} \implies \text{of-complex } x \in \mathbb{Z}$
 ⟨proof⟩

Every complex algebra has characteristic zero.

lemma *fraction-scaleC-times* [simp]:
 fixes $a :: 'a::\text{complex-algebra-1}$
 shows $(\text{numeral } u / \text{numeral } v) *_{\mathbb{C}} (\text{numeral } w * a) = (\text{numeral } u * \text{numeral } w / \text{numeral } v) *_{\mathbb{C}} a$
 ⟨proof⟩

lemma *inverse-scaleC-times* [simp]:
 fixes $a :: 'a::\text{complex-algebra-1}$
 shows $(1 / \text{numeral } v) *_{\mathbb{C}} (\text{numeral } w * a) = (\text{numeral } w / \text{numeral } v) *_{\mathbb{C}} a$
 ⟨proof⟩

lemma *scaleC-times* [simp]:
 fixes $a :: 'a::\text{complex-algebra-1}$
 shows $(\text{numeral } u) *_{\mathbb{C}} (\text{numeral } w * a) = (\text{numeral } u * \text{numeral } w) *_{\mathbb{C}} a$
 ⟨proof⟩

6.3 The Set of Real Numbers

definition *Complexs* :: $'a::\text{complex-algebra-1}$ set ($\mathbb{C}$)
 where $\mathbb{C} = \text{range of-complex}$

lemma *Complexs-of-complex* [simp]: $\text{of-complex } r \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-of-int* [simp]: $\text{of-int } z \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-of-nat* [simp]: $\text{of-nat } n \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-numeral* [simp]: $\text{numeral } w \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-0* [simp]: $0 \in \mathbb{C}$ and *Complexs-1* [simp]: $1 \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-add* [simp]: $a \in \mathbb{C} \implies b \in \mathbb{C} \implies a + b \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-minus* [simp]: $a \in \mathbb{C} \implies -a \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-minus-iff* [simp]: $-a \in \mathbb{C} \longleftrightarrow a \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-diff* [simp]: $a \in \mathbb{C} \implies b \in \mathbb{C} \implies a - b \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-mult* [simp]: $a \in \mathbb{C} \implies b \in \mathbb{C} \implies a * b \in \mathbb{C}$
 ⟨proof⟩

lemma *nonzero-Complexs-inverse*: $a \in \mathbb{C} \implies a \neq 0 \implies \text{inverse } a \in \mathbb{C}$
 for $a :: 'a :: \text{complex-div-algebra}$
 ⟨proof⟩

lemma *Complexs-inverse*: $a \in \mathbb{C} \implies \text{inverse } a \in \mathbb{C}$
 for $a :: 'a :: \{\text{complex-div-algebra}, \text{division-ring}\}$
 ⟨proof⟩

lemma *Complexs-inverse-iff* [simp]: $\text{inverse } x \in \mathbb{C} \longleftrightarrow x \in \mathbb{C}$
 for $x :: 'a :: \{\text{complex-div-algebra}, \text{division-ring}\}$
 ⟨proof⟩

lemma *nonzero-Complexs-divide*: $a \in \mathbb{C} \implies b \in \mathbb{C} \implies b \neq 0 \implies a / b \in \mathbb{C}$
 for $a \ b :: 'a :: \text{complex-field}$
 ⟨proof⟩

lemma *Complexs-divide* [simp]: $a \in \mathbb{C} \implies b \in \mathbb{C} \implies a / b \in \mathbb{C}$
 for $a \ b :: 'a :: \{\text{complex-field}, \text{field}\}$
 ⟨proof⟩

lemma *Complexs-power* [simp]: $a \in \mathbb{C} \implies a ^ n \in \mathbb{C}$
 for $a :: 'a :: \text{complex-algebra-1}$
 ⟨proof⟩

lemma *Complexs-cases* [cases set: *Complexs*]:
 assumes $q \in \mathbb{C}$
 obtains (of-complex) c where $q = \text{of-complex } c$
 ⟨proof⟩

lemma *sum-in-Complexs* [intro,simp]: $(\bigwedge i. i \in s \implies f \ i \in \mathbb{C}) \implies \text{sum } f \ s \in \mathbb{C}$
 ⟨proof⟩

lemma *prod-in-Complexs* [intro,simp]: $(\bigwedge i. i \in s \implies f \ i \in \mathbb{C}) \implies \text{prod } f \ s \in \mathbb{C}$
 ⟨proof⟩

lemma *Complexs-induct* [case-names of-complex, induct set: *Complexs*]:
 $q \in \mathbb{C} \implies (\bigwedge r. P \ (\text{of-complex } r)) \implies P \ q$
 ⟨proof⟩

6.4 Ordered complex vector spaces

class *ordered-complex-vector* = *complex-vector* + *ordered-ab-group-add* +
 assumes *scaleC-left-mono*: $x \leq y \implies 0 \leq a \implies a *_C x \leq a *_C y$

and *scaleC-right-mono*: $a \leq b \implies 0 \leq x \implies a *_C x \leq b *_C x$
begin

subclass (**in** *ordered-complex-vector*) *ordered-real-vector*
 $\langle proof \rangle$

lemma *scaleC-mono*:
 $a \leq b \implies x \leq y \implies 0 \leq b \implies 0 \leq x \implies a *_C x \leq b *_C y$
 $\langle proof \rangle$

lemma *scaleC-mono'*:
 $a \leq b \implies c \leq d \implies 0 \leq a \implies 0 \leq c \implies a *_C c \leq b *_C d$
 $\langle proof \rangle$

lemma *pos-le-divideC-eq* [*field-simps*]:
 $a \leq b /_C c \longleftrightarrow c *_C a \leq b$ (**is** $?P \longleftrightarrow ?Q$) **if** $0 < c$
 $\langle proof \rangle$

lemma *pos-less-divideC-eq* [*field-simps*]:
 $a < b /_C c \longleftrightarrow c *_C a < b$ **if** $c > 0$
 $\langle proof \rangle$

lemma *pos-divideC-le-eq* [*field-simps*]:
 $b /_C c \leq a \longleftrightarrow b \leq c *_C a$ **if** $c > 0$
 $\langle proof \rangle$

lemma *pos-divideC-less-eq* [*field-simps*]:
 $b /_C c < a \longleftrightarrow b < c *_C a$ **if** $c > 0$
 $\langle proof \rangle$

lemma *pos-le-minus-divideC-eq* [*field-simps*]:
 $a \leq - (b /_C c) \longleftrightarrow c *_C a \leq - b$ **if** $c > 0$
 $\langle proof \rangle$

lemma *pos-less-minus-divideC-eq* [*field-simps*]:
 $a < - (b /_C c) \longleftrightarrow c *_C a < - b$ **if** $c > 0$
 $\langle proof \rangle$

lemma *pos-minus-divideC-le-eq* [*field-simps*]:
 $- (b /_C c) \leq a \longleftrightarrow - b \leq c *_C a$ **if** $c > 0$
 $\langle proof \rangle$

lemma *pos-minus-divideC-less-eq* [*field-simps*]:
 $- (b /_C c) < a \longleftrightarrow - b < c *_C a$ **if** $c > 0$
 $\langle proof \rangle$

lemma *scaleC-image-atLeastAtMost*: $c > 0 \implies scaleC\ c\ ' \{x..y\} = \{c *_C x..c *_C y\}$
 $\langle proof \rangle$

end

lemma *neg-le-divideC-eq* [*field-simps*]:
 $a \leq b /_C c \longleftrightarrow b \leq c *_C a$ (**is** $?P \longleftrightarrow ?Q$) **if** $c < 0$
for $a\ b :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *neg-less-divideC-eq* [*field-simps*]:
 $a < b /_C c \longleftrightarrow b < c *_C a$ **if** $c < 0$
for $a\ b :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *neg-divideC-le-eq* [*field-simps*]:
 $b /_C c \leq a \longleftrightarrow c *_C a \leq b$ **if** $c < 0$
for $a\ b :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *neg-divideC-less-eq* [*field-simps*]:
 $b /_C c < a \longleftrightarrow c *_C a < b$ **if** $c < 0$
for $a\ b :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *neg-le-minus-divideC-eq* [*field-simps*]:
 $a \leq - (b /_C c) \longleftrightarrow - b \leq c *_C a$ **if** $c < 0$
for $a\ b :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *neg-less-minus-divideC-eq* [*field-simps*]:
 $a < - (b /_C c) \longleftrightarrow - b < c *_C a$ **if** $c < 0$
for $a\ b :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *neg-minus-divideC-le-eq* [*field-simps*]:
 $-(b /_C c) \leq a \longleftrightarrow c *_C a \leq -b$ **if** $c < 0$
for $a\ b :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *neg-minus-divideC-less-eq* [*field-simps*]:
 $-(b /_C c) < a \longleftrightarrow c *_C a < -b$ **if** $c < 0$
for $a\ b :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *divideC-field-splits-simps-1* [*field-split-simps*]:
 $a = b /_C c \longleftrightarrow (\text{if } c = 0 \text{ then } a = 0 \text{ else } c *_C a = b)$
 $b /_C c = a \longleftrightarrow (\text{if } c = 0 \text{ then } a = 0 \text{ else } b = c *_C a)$
 $a + b /_C c = (\text{if } c = 0 \text{ then } a \text{ else } (c *_C a + b) /_C c)$
 $a /_C c + b = (\text{if } c = 0 \text{ then } b \text{ else } (a + c *_C b) /_C c)$
 $a - b /_C c = (\text{if } c = 0 \text{ then } a \text{ else } (c *_C a - b) /_C c)$

$a /_C c - b = (\text{if } c = 0 \text{ then } -b \text{ else } (a - c *_C b) /_C c)$
 $-(a /_C c) + b = (\text{if } c = 0 \text{ then } b \text{ else } (-a + c *_C b) /_C c)$
 $-(a /_C c) - b = (\text{if } c = 0 \text{ then } -b \text{ else } (-a - c *_C b) /_C c)$
for $a \ b :: 'a :: \text{complex-vector}$
 $\langle \text{proof} \rangle$

lemma *divideC-field-splits-simps-2* [*field-split-simps*]:

$0 < c \implies a \leq b /_C c \longleftrightarrow (\text{if } c > 0 \text{ then } c *_C a \leq b \text{ else if } c < 0 \text{ then } b \leq c *_C a \text{ else } a \leq 0)$
 $0 < c \implies a < b /_C c \longleftrightarrow (\text{if } c > 0 \text{ then } c *_C a < b \text{ else if } c < 0 \text{ then } b < c *_C a \text{ else } a < 0)$
 $0 < c \implies b /_C c \leq a \longleftrightarrow (\text{if } c > 0 \text{ then } b \leq c *_C a \text{ else if } c < 0 \text{ then } c *_C a \leq b \text{ else } a \geq 0)$
 $0 < c \implies b /_C c < a \longleftrightarrow (\text{if } c > 0 \text{ then } b < c *_C a \text{ else if } c < 0 \text{ then } c *_C a < b \text{ else } a > 0)$
 $0 < c \implies a \leq -(b /_C c) \longleftrightarrow (\text{if } c > 0 \text{ then } c *_C a \leq -b \text{ else if } c < 0 \text{ then } -b \leq c *_C a \text{ else } a \leq 0)$
 $0 < c \implies a < -(b /_C c) \longleftrightarrow (\text{if } c > 0 \text{ then } c *_C a < -b \text{ else if } c < 0 \text{ then } -b < c *_C a \text{ else } a < 0)$
 $0 < c \implies -(b /_C c) \leq a \longleftrightarrow (\text{if } c > 0 \text{ then } -b \leq c *_C a \text{ else if } c < 0 \text{ then } c *_C a \leq -b \text{ else } a \geq 0)$
 $0 < c \implies -(b /_C c) < a \longleftrightarrow (\text{if } c > 0 \text{ then } -b < c *_C a \text{ else if } c < 0 \text{ then } c *_C a < -b \text{ else } a > 0)$
for $a \ b :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-nonneg-nonneg*: $0 \leq a \implies 0 \leq x \implies 0 \leq a *_C x$

for $x :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-nonneg-nonpos*: $0 \leq a \implies x \leq 0 \implies a *_C x \leq 0$

for $x :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-nonpos-nonneg*: $a \leq 0 \implies 0 \leq x \implies a *_C x \leq 0$

for $x :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *split-scaleC-neg-le*: $(0 \leq a \wedge x \leq 0) \vee (a \leq 0 \wedge 0 \leq x) \implies a *_C x \leq 0$

for $x :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *cle-add-iff1*: $a *_C e + c \leq b *_C e + d \longleftrightarrow (a - b) *_C e + c \leq d$

for $c \ d \ e :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *cle-add-iff2*: $a *_C e + c \leq b *_C e + d \longleftrightarrow c \leq (b - a) *_C e + d$

for $c \ d \ e :: 'a :: \text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-left-mono-neg*: $b \leq a \implies c \leq 0 \implies c *_C a \leq c *_C b$
for $a \ b :: 'a::\text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-right-mono-neg*: $b \leq a \implies c \leq 0 \implies a *_C c \leq b *_C c$
for $c :: 'a::\text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-nonpos-nonpos*: $a \leq 0 \implies b \leq 0 \implies 0 \leq a *_C b$
for $b :: 'a::\text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *split-scaleC-pos-le*: $(0 \leq a \wedge 0 \leq b) \vee (a \leq 0 \wedge b \leq 0) \implies 0 \leq a *_C b$
for $b :: 'a::\text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *zero-le-scaleC-iff*:
fixes $b :: 'a::\text{ordered-complex-vector}$
assumes $a \in \mathbb{R}$
shows $0 \leq a *_C b \longleftrightarrow 0 < a \wedge 0 \leq b \vee a < 0 \wedge b \leq 0 \vee a = 0$
(is ?lhs = ?rhs)
 $\langle \text{proof} \rangle$

lemma *scaleC-le-0-iff*:
 $a *_C b \leq 0 \longleftrightarrow 0 < a \wedge b \leq 0 \vee a < 0 \wedge 0 \leq b \vee a = 0$
if $a \in \mathbb{R}$
for $b :: 'a::\text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-le-cancel-left*: $c *_C a \leq c *_C b \longleftrightarrow (0 < c \longrightarrow a \leq b) \wedge (c < 0 \longrightarrow b \leq a)$
if $c \in \mathbb{R}$
for $b :: 'a::\text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-le-cancel-left-pos*: $0 < c \implies c *_C a \leq c *_C b \longleftrightarrow a \leq b$
for $b :: 'a::\text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-le-cancel-left-neg*: $c < 0 \implies c *_C a \leq c *_C b \longleftrightarrow b \leq a$
for $b :: 'a::\text{ordered-complex-vector}$
 $\langle \text{proof} \rangle$

lemma *scaleC-left-le-one-le*: $0 \leq x \implies a \leq 1 \implies a *_C x \leq x$
for $x :: 'a::\text{ordered-complex-vector}$ **and** $a :: \text{complex}$
 $\langle \text{proof} \rangle$

6.5 Complex normed vector spaces

```
class complex-normed-vector = complex-vector + sgn-div-norm + dist-norm +
uniformity-dist + open-uniformity +
  real-normed-vector +
  assumes norm-scaleC [simp]: norm (scaleC a x) = cmod a * norm x
begin
```

```
end
```

```
class complex-normed-algebra = complex-algebra + complex-normed-vector +
  real-normed-algebra
```

```
class complex-normed-algebra-1 = complex-algebra-1 + complex-normed-algebra +
  real-normed-algebra-1
```

```
lemma (in complex-normed-algebra-1) scaleC-power [simp]: (scaleC x y) ^ n =
scaleC (x ^ n) (y ^ n)
  ⟨proof⟩
```

```
class complex-normed-div-algebra = complex-div-algebra + complex-normed-vector
+
  real-normed-div-algebra
```

```
class complex-normed-field = complex-field + complex-normed-div-algebra
```

```
subclass (in complex-normed-field) real-normed-field ⟨proof⟩
```

```
instance complex-normed-div-algebra < complex-normed-algebra-1 ⟨proof⟩
```

```
context complex-normed-vector begin
```

```
end
```

```
lemma dist-scaleC [simp]: dist (x *C a) (y *C a) = |x - y| * norm a
  for a :: 'a::complex-normed-vector
  ⟨proof⟩
```

```
lemma norm-of-complex [simp]: norm (of-complex c :: 'a::complex-normed-algebra-1)
= cmod c
  ⟨proof⟩
```

lemma *norm-of-complex-add1* [simp]: *norm (of-complex x + 1 :: 'a :: complex-normed-div-algebra)*
 $= \text{cmod } (x + 1)$
 ⟨proof⟩

lemma *norm-of-complex-addn* [simp]:
 $\text{norm } (\text{of-complex } x + \text{numeral } b :: 'a :: \text{complex-normed-div-algebra}) = \text{cmod } (x + \text{numeral } b)$
 ⟨proof⟩

lemma *norm-of-complex-diff* [simp]:
 $\text{norm } (\text{of-complex } b - \text{of-complex } a :: 'a :: \text{complex-normed-div-algebra}) \leq \text{cmod } (b - a)$
 ⟨proof⟩

6.6 Metric spaces

Every normed vector space is a metric space.

6.7 Class instances for complex numbers

instantiation *complex* :: *complex-normed-field*
begin

instance
 ⟨proof⟩

end

declare *uniformity-Absort*[**where** 'a=complex, code]

lemma *dist-of-complex* [simp]: $\text{dist } (\text{of-complex } x :: 'a) (\text{of-complex } y) = \text{dist } x y$
for *a* :: 'a :: complex-normed-div-algebra
 ⟨proof⟩

declare [[code abort: open :: complex set $\Rightarrow$ bool]]

lemma *closed-complex-atMost*: $\langle \text{closed } \{..a :: \text{complex}\} \rangle$
 ⟨proof⟩

lemma *closed-complex-atLeast*: $\langle \text{closed } \{a :: \text{complex}..\} \rangle$
 ⟨proof⟩

lemma *closed-complex-atLeastAtMost*: $\langle \text{closed } \{a :: \text{complex} .. b\} \rangle$
 ⟨proof⟩

6.8 Sign function

lemma *sgn-scaleC*: $\text{sgn} (\text{scaleC } r \ x) = \text{scaleC } (\text{sgn } r) (\text{sgn } x)$
for $x :: 'a::\text{complex-normed-vector}$
 $\langle \text{proof} \rangle$

lemma *sgn-of-complex*: $\text{sgn} (\text{of-complex } r :: 'a::\text{complex-normed-algebra-1}) = \text{of-complex} (\text{sgn } r)$
 $\langle \text{proof} \rangle$

lemma *complex-sgn-eq*: $\text{sgn } x = x / |x|$
for $x :: \text{complex}$
 $\langle \text{proof} \rangle$

lemma *czero-le-sgn-iff* [simp]: $0 \leq \text{sgn } x \longleftrightarrow 0 \leq x$
for $x :: \text{complex}$
 $\langle \text{proof} \rangle$

lemma *csgn-le-0-iff* [simp]: $\text{sgn } x \leq 0 \longleftrightarrow x \leq 0$
for $x :: \text{complex}$
 $\langle \text{proof} \rangle$

6.9 Bounded Linear and Bilinear Operators

lemma *clinearI*: *clinear* f
if $\bigwedge b1 \ b2. f (b1 + b2) = f \ b1 + f \ b2$
 $\bigwedge r \ b. f (r *_C b) = r *_C f \ b$
 $\langle \text{proof} \rangle$

lemma *clinear-iff*:
 $\text{clinear } f \longleftrightarrow (\forall x \ y. f (x + y) = f \ x + f \ y) \wedge (\forall c \ x. f (c *_C x) = c *_C f \ x)$
(is *clinear* $f \longleftrightarrow ?rhs)$
 $\langle \text{proof} \rangle$

lemmas *clinear-scaleC-left* = *complex-vector.linear-scale-left*
lemmas *clinear-imp-scaleC* = *complex-vector.linear-imp-scale*

corollary *complex-clinearD*:
fixes $f :: \text{complex} \Rightarrow \text{complex}$
assumes *clinear* f **obtains** c **where** $f = (*) \ c$
 $\langle \text{proof} \rangle$

lemma *clinear-times-of-complex*: *clinear* $(\lambda x. a * \text{of-complex } x)$
 $\langle \text{proof} \rangle$

locale *bounded-clinear* = *clinear* f **for** $f :: 'a::\text{complex-normed-vector} \Rightarrow 'b::\text{complex-normed-vector}$
+
assumes *bounded*: $\exists K. \forall x. \text{norm } (f \ x) \leq \text{norm } x * K$
begin

sublocale *real*: *bounded-linear*

— Gives access to all lemmas from *bounded-linear* using prefix *real*.

<proof>

lemmas *pos-bounded* = *real.pos-bounded*

lemmas *nonneg-bounded* = *real.nonneg-bounded*

lemma *clinear*: *clinear* *f*

<proof>

end

lemma *bounded-clinear-intro*:

assumes $\bigwedge x y. f (x + y) = f x + f y$

and $\bigwedge r x. f (scaleC\ r\ x) = scaleC\ r\ (f\ x)$

and $\bigwedge x. norm\ (f\ x) \leq norm\ x * K$

shows *bounded-clinear* *f*

<proof>

locale *bounded-cbilinear* =

fixes *prod* :: '*a*::*complex-normed-vector* $\Rightarrow$ '*b*::*complex-normed-vector* $\Rightarrow$ '*c*::*complex-normed-vector*
(**infixl** $\langle ** \rangle$ 70)

assumes *add-left*: *prod* (*a* + *a'*) *b* = *prod* *a* *b* + *prod* *a'* *b*

and *add-right*: *prod* *a* (*b* + *b'*) = *prod* *a* *b* + *prod* *a* *b'*

and *scaleC-left*: *prod* (*scaleC* *r* *a*) *b* = *scaleC* *r* (*prod* *a* *b*)

and *scaleC-right*: *prod* *a* (*scaleC* *r* *b*) = *scaleC* *r* (*prod* *a* *b*)

and *bounded*: $\exists K. \forall a\ b. norm\ (prod\ a\ b) \leq norm\ a * norm\ b * K$

begin

sublocale *real*: *bounded-bilinear*

— Gives access to all lemmas from *bounded-bilinear* using prefix *real*.

<proof>

lemmas *pos-bounded* = *real.pos-bounded*

lemmas *nonneg-bounded* = *real.nonneg-bounded*

lemmas *additive-right* = *real.additive-right*

lemmas *additive-left* = *real.additive-left*

lemmas *zero-left* = *real.zero-left*

lemmas *zero-right* = *real.zero-right*

lemmas *minus-left* = *real.minus-left*

lemmas *minus-right* = *real.minus-right*

lemmas *diff-left* = *real.diff-left*

lemmas *diff-right* = *real.diff-right*

lemmas *sum-left* = *real.sum-left*

lemmas *sum-right* = *real.sum-right*

lemmas *prod-diff-prod* = *real.prod-diff-prod*

lemma *bounded-clinear-left*: *bounded-clinear* ($\lambda a. a ** b$)
<proof>

lemma *bounded-clinear-right*: *bounded-clinear* ($\lambda b. a ** b$)
<proof>

lemma *flip*: *bounded-cbilinear* ($\lambda x y. y ** x$)
<proof>

lemma *comp1*:
 assumes *bounded-clinear* *g*
 shows *bounded-cbilinear* ($\lambda x. (**) (g x)$)
<proof>

lemma *comp*: *bounded-clinear* *f* $\implies$ *bounded-clinear* *g* $\implies$ *bounded-cbilinear* (λx
*y. f x ** g y*)
<proof>

end

lemma *bounded-clinear-ident[simp]*: *bounded-clinear* ($\lambda x. x$)
<proof>

lemma *bounded-clinear-zero[simp]*: *bounded-clinear* ($\lambda x. 0$)
<proof>

lemma *bounded-clinear-add*:
 assumes *bounded-clinear* *f*
 and *bounded-clinear* *g*
 shows *bounded-clinear* ($\lambda x. f x + g x$)
<proof>

lemma *bounded-clinear-minus*:
 assumes *bounded-clinear* *f*
 shows *bounded-clinear* ($\lambda x. - f x$)
<proof>

lemma *bounded-clinear-sub*: *bounded-clinear* *f* $\implies$ *bounded-clinear* *g* $\implies$ *bounded-clinear*
($\lambda x. f x - g x$)
<proof>

lemma *bounded-clinear-sum*:
 fixes *f* :: '*i* $\Rightarrow$ 'a::complex-normed-vector $\Rightarrow$ 'b::complex-normed-vector
 shows ($\bigwedge i. i \in I \implies \text{bounded-clinear } (f i) \implies \text{bounded-clinear } (\lambda x. \sum_{i \in I. f$
i x))
<proof>

lemma *bounded-clinear-compose*:
assumes *bounded-clinear* f
and *bounded-clinear* g
shows *bounded-clinear* $(\lambda x. f (g x))$
 $\langle proof \rangle$

lemma *bounded-cbilinear-mult*: *bounded-cbilinear* $((*) :: 'a \Rightarrow 'a \Rightarrow 'a :: \text{complex-normed-algebra})$
 $\langle proof \rangle$

lemma *bounded-clinear-mult-left*: *bounded-clinear* $(\lambda x :: 'a :: \text{complex-normed-algebra}. x * y)$
 $\langle proof \rangle$

lemma *bounded-clinear-mult-right*: *bounded-clinear* $(\lambda y :: 'a :: \text{complex-normed-algebra}. x * y)$
 $\langle proof \rangle$

lemmas *bounded-clinear-mult-const* =
bounded-clinear-mult-left [THEN *bounded-clinear-compose*]

lemmas *bounded-clinear-const-mult* =
bounded-clinear-mult-right [THEN *bounded-clinear-compose*]

lemma *bounded-clinear-divide*: *bounded-clinear* $(\lambda x. x / y)$
for $y :: 'a :: \text{complex-normed-field}$
 $\langle proof \rangle$

lemma *bounded-cbilinear-scaleC*: *bounded-cbilinear* *scaleC*
 $\langle proof \rangle$

lemma *bounded-clinear-scaleC-left*: *bounded-clinear* $(\lambda c. \text{scaleC } c \ x)$
 $\langle proof \rangle$

lemma *bounded-clinear-scaleC-right*: *bounded-clinear* $(\lambda x. \text{scaleC } c \ x)$
 $\langle proof \rangle$

lemmas *bounded-clinear-scaleC-const* =
bounded-clinear-scaleC-left [THEN *bounded-clinear-compose*]

lemmas *bounded-clinear-const-scaleC* =
bounded-clinear-scaleC-right [THEN *bounded-clinear-compose*]

lemma *bounded-clinear-of-complex*: *bounded-clinear* $(\lambda r. \text{of-complex } r)$
 $\langle proof \rangle$

lemma *complex-bounded-clinear*: *bounded-clinear* $f \longleftrightarrow (\exists c :: \text{complex}. f = (\lambda x. x * c))$
for $f :: \text{complex} \Rightarrow \text{complex}$
 $\langle proof \rangle$

6.9.1 Limits of Sequences

6.10 Cauchy sequences

lemma *cCauchy-iff2*: $\text{Cauchy } X \longleftrightarrow (\forall j. (\exists M. \forall m \geq M. \forall n \geq M. \text{cmod } (X\ m - X\ n) < \text{inverse } (\text{real } (\text{Suc } j))))$
 $\langle \text{proof} \rangle$

6.11 The set of complex numbers is a complete metric space

Proof that Cauchy sequences converge based on the one from <http://pirate.shu.edu/~wachsmut/ira/numseq/proofs/cauconv.html>

If sequence X is Cauchy, then its limit is the lub of $\{r. \exists N. \forall n \geq N. r < X\ n\}$

lemma *complex-increasing-LIMSEQ*:
fixes $f :: \text{nat} \Rightarrow \text{complex}$
assumes $\text{inc}: \bigwedge n. f\ n \leq f\ (\text{Suc } n)$
and $\text{bdd}: \bigwedge n. f\ n \leq l$
and $\text{en}: \bigwedge e. 0 < e \implies \exists n. l \leq f\ n + e$
shows $f \longrightarrow l$
 $\langle \text{proof} \rangle$

lemma *complex-Cauchy-convergent*:
fixes $X :: \text{nat} \Rightarrow \text{complex}$
assumes $X: \text{Cauchy } X$
shows $\text{convergent } X$
 $\langle \text{proof} \rangle$

instance *complex* :: *complete-space*
 $\langle \text{proof} \rangle$

class *cbanach* = *complex-normed-vector* + *complete-space*

subclass (**in** *cbanach*) *banach* $\langle \text{proof} \rangle$

instance *complex* :: *banach* $\langle \text{proof} \rangle$

end

7 Complex-Vector-Spaces – Complex Vector Spaces

theory *Complex-Vector-Spaces*
imports

HOL–Analysis.Elementary-Topology
HOL–Analysis.Operator-Norm
HOL–Analysis.Elementary-Normed-Spaces
HOL–Library.Set-Algebras
HOL–Analysis.Starlike
HOL–Types-To-Sets.Types-To-Sets
HOL–Library.Complemented-Lattices
HOL–Library.Function-Algebras

Extra-Vector-Spaces
Extra-Ordered-Fields
Extra-Operator-Norm
Extra-General

Complex-Vector-Spaces0

begin

bundle *norm-syntax* **begin**

notation *norm* ($\langle \cdot \| \cdot \rangle$)

end

unbundle *lattice-syntax*

7.1 Misc

lemma (**in** *vector-space*) *span-image-scale'*:

— Strengthening of *vector-space.span-image-scale* without the condition *finite S*

assumes *nz*: $\bigwedge x. x \in S \implies c \ x \neq 0$

shows *span* $((\lambda x. c \ x * s \ x) \ ` S) = \text{span } S$

$\langle \text{proof} \rangle$

lemma (**in** *scaleC*) *scaleC-real*: **assumes** $r \in \mathbb{R}$ **shows** $r *_{\mathbb{C}} x = \text{Re } r *_{\mathbb{R}} x$

$\langle \text{proof} \rangle$

lemma *of-complex-of-real-eq* [*simp*]: *of-complex* (*of-real* n) = *of-real* n

$\langle \text{proof} \rangle$

lemma *Complexs-of-real* [*simp*]: *of-real* $r \in \mathbb{C}$

$\langle \text{proof} \rangle$

lemma *Reals-in-Complexs*: $\mathbb{R} \subseteq \mathbb{C}$

$\langle \text{proof} \rangle$

lemma (**in** *bounded-clinear*) *bounded-linear*: *bounded-linear* f

$\langle \text{proof} \rangle$

lemma *clinear-times*: *clinear* $(\lambda x. c * x)$

for $c :: 'a :: \text{complex-algebra}$

$\langle \text{proof} \rangle$

lemma (in *clinear*) *linear*: $\langle \text{linear } f \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-clinearI*:
assumes $\langle \bigwedge b1\ b2. f\ (b1 + b2) = f\ b1 + f\ b2 \rangle$
assumes $\langle \bigwedge r\ b. f\ (r *_C\ b) = r *_C\ f\ b \rangle$
assumes $\langle \bigwedge x. \text{norm}\ (f\ x) \leq \text{norm}\ x * K \rangle$
shows *bounded-clinear* *f*
 $\langle \text{proof} \rangle$

lemma *bounded-clinear-id[simp]*: $\langle \text{bounded-clinear id} \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-clinear-0[simp]*: $\langle \text{bounded-clinear } 0 \rangle$
 $\langle \text{proof} \rangle$

definition *cbilinear* :: $\langle 'a::\text{complex-vector} \Rightarrow 'b::\text{complex-vector} \Rightarrow 'c::\text{complex-vector} \rangle$
 $\Rightarrow \text{bool}$
where $\langle \text{cbilinear} = (\lambda f. (\forall y. \text{clinear}\ (\lambda x. f\ x\ y)) \wedge (\forall x. \text{clinear}\ (\lambda y. f\ x\ y))) \rangle$

lemma *cbilinear-add-left*:
assumes $\langle \text{cbilinear } f \rangle$
shows $\langle f\ (a + b)\ c = f\ a\ c + f\ b\ c \rangle$
 $\langle \text{proof} \rangle$

lemma *cbilinear-add-right*:
assumes $\langle \text{cbilinear } f \rangle$
shows $\langle f\ a\ (b + c) = f\ a\ b + f\ a\ c \rangle$
 $\langle \text{proof} \rangle$

lemma *cbilinear-times*:
fixes $g' :: \langle 'a::\text{complex-vector} \Rightarrow \text{complex} \rangle$ **and** $g :: \langle 'b::\text{complex-vector} \Rightarrow \text{complex} \rangle$
assumes $\langle \bigwedge x\ y. h\ x\ y = (g'\ x) * (g\ y) \rangle$ **and** $\langle \text{clinear } g \rangle$ **and** $\langle \text{clinear } g' \rangle$
shows $\langle \text{cbilinear } h \rangle$
 $\langle \text{proof} \rangle$

lemma *csubspace-is-subspace*: $\text{csubspace } A \Longrightarrow \text{subspace } A$
 $\langle \text{proof} \rangle$

lemma *span-subset-cspan*: $\text{span } A \subseteq \text{cspan } A$
 $\langle \text{proof} \rangle$

lemma *cinddependent-implies-independent*:
assumes *cinddependent* $(S :: 'a::\text{complex-vector set})$
shows *independent* *S*

⟨proof⟩

lemma *cspan-singleton*: $cspan \{x\} = \{\alpha *_{\mathcal{C}} x \mid \alpha. True\}$
 ⟨proof⟩

lemma *cspan-as-span*:
 $cspan (B :: 'a :: complex-vector set) = span (B \cup scale_{\mathcal{C}} i \text{ ` } B)$
 ⟨proof⟩

lemma *isomorphic-equal-cdim*:
 assumes *lin-f*: ⟨linear f ⟩
 assumes *inj-f*: ⟨inj-on f ($cspan S$)⟩
 assumes *im-S*: ⟨ $f \text{ ` } S = T$ ⟩
 shows ⟨ $cdim S = cdim T$ ⟩
 ⟨proof⟩

lemma *cindependent-inter-scaleC-cindependent*:
 assumes *a1*: *cindependent* ($B :: 'a :: complex-vector set$) and *a3*: $c \neq 1$
 shows $B \cap (*_{\mathcal{C}}) c \text{ ` } B = \{\}$
 ⟨proof⟩

lemma *real-independent-from-complex-independent*:
 assumes *cindependent* ($B :: 'a :: complex-vector set$)
 defines $B' == ((*_{\mathcal{C}}) i \text{ ` } B)$
 shows *independent* ($B \cup B'$)
 ⟨proof⟩

lemma *crepresentation-from-representation*:
 assumes *a1*: *cindependent* B and *a2*: $b \in B$ and *a3*: *finite* B
 shows $crepresentation B \psi b = (representation (B \cup (*_{\mathcal{C}}) i \text{ ` } B) \psi b) + i *_{\mathcal{C}} (representation (B \cup (*_{\mathcal{C}}) i \text{ ` } B) \psi (i *_{\mathcal{C}} b))$
 ⟨proof⟩

lemma *CARD-1-vec-0[simp]*: ⟨ $(\psi :: - :: \{complex-vector, CARD-1\}) = 0$ ⟩
 ⟨proof⟩

lemma *scaleC-cindependent*:
 assumes *a1*: *cindependent* ($B :: 'a :: complex-vector set$) and *a3*: $c \neq 0$
 shows *cindependent* ($(*_{\mathcal{C}}) c \text{ ` } B$)
 ⟨proof⟩

lemma *cspan-eqI*:
 assumes ⟨ $\bigwedge a. a \in A \implies a \in cspan B$ ⟩
 assumes ⟨ $\bigwedge b. b \in B \implies b \in cspan A$ ⟩

shows $\langle \text{cspan } A = \text{cspan } B \rangle$
 $\langle \text{proof} \rangle$

lemma (in bounded-cbilinear) bounded-bilinear[simp]: bounded-bilinear prod
 $\langle \text{proof} \rangle$

lemma norm-scaleC-sgn[simp]: $\langle \text{complex-of-real } (\text{norm } \psi) *_{\mathbb{C}} \text{sgn } \psi = \psi \rangle$ **for** ψ
 $:: 'a::\text{complex-normed-vector}$
 $\langle \text{proof} \rangle$

lemma scaleC-of-complex[simp]: $\langle \text{scaleC } x (\text{of-complex } y) = \text{of-complex } (x * y) \rangle$
 $\langle \text{proof} \rangle$

lemma bounded-clinear-inv:
assumes [simp]: $\langle \text{bounded-clinear } f \rangle$
assumes b: $\langle b > 0 \rangle$
assumes bound: $\langle \bigwedge x. \text{norm } (f x) \geq b * \text{norm } x \rangle$
assumes $\langle \text{surj } f \rangle$
shows $\langle \text{bounded-clinear } (\text{inv } f) \rangle$
 $\langle \text{proof} \rangle$

lemma range-is-csubspace[simp]:
assumes a1: $\langle \text{clinear } f \rangle$
shows $\langle \text{csubspace } (\text{range } f) \rangle$
 $\langle \text{proof} \rangle$

lemma csubspace-is-convex[simp]:
assumes a1: $\langle \text{csubspace } M \rangle$
shows $\langle \text{convex } M \rangle$
 $\langle \text{proof} \rangle$

lemma kernel-is-csubspace[simp]:
assumes a1: $\langle \text{clinear } f \rangle$
shows $\langle \text{csubspace } (f^{-1} \{0\}) \rangle$
 $\langle \text{proof} \rangle$

lemma bounded-cbilinear-0[simp]: $\langle \text{bounded-cbilinear } (\lambda -. 0) \rangle$
 $\langle \text{proof} \rangle$

lemma bounded-cbilinear-0'[simp]: $\langle \text{bounded-cbilinear } 0 \rangle$
 $\langle \text{proof} \rangle$

lemma bounded-cbilinear-apply-bounded-clinear: $\langle \text{bounded-clinear } (f x) \rangle$ **if** $\langle \text{bounded-cbilinear } f \rangle$
 $\langle \text{proof} \rangle$

lemma clinear-scaleR[simp]: $\langle \text{clinear } (\text{scaleR } x) \rangle$
 $\langle \text{proof} \rangle$

lemma *abs-summable-on-scaleC-left* [intro]:
fixes $c :: \langle 'a :: \text{complex-normed-vector} \rangle$
assumes $c \neq 0 \implies f \text{ abs-summable-on } A$
shows $(\lambda x. f\ x *_{\mathbb{C}} c) \text{ abs-summable-on } A$
 $\langle \text{proof} \rangle$

lemma *abs-summable-on-scaleC-right* [intro]:
fixes $f :: \langle 'a \Rightarrow 'b :: \text{complex-normed-vector} \rangle$
assumes $c \neq 0 \implies f \text{ abs-summable-on } A$
shows $(\lambda x. c *_{\mathbb{C}} f\ x) \text{ abs-summable-on } A$
 $\langle \text{proof} \rangle$

lemma *clinear-of-complex*[iff]: $\langle \text{clinear of-complex} \rangle$
 $\langle \text{proof} \rangle$

lemma *infsun-of-bool-scaleC*: $\langle (\sum_{x \in X}. \text{of-bool } (x=y) *_{\mathbb{C}} f\ x) = \text{of-bool } (y \in X) *_{\mathbb{C}} f\ y \rangle$ **for** $f :: \langle - \Rightarrow - :: \text{complex-vector} \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-bounded-clinear*:
assumes *bounded-clinear* h **and** $(f \text{ has-sum } S) \ A$
shows $((\lambda x. h\ (f\ x)) \text{ has-sum } h\ S) \ A$
 $\langle \text{proof} \rangle$

lemma *scaleC-scaleR-commute*: $\langle a *_{\mathbb{C}} b *_{\mathbb{R}} x = b *_{\mathbb{R}} a *_{\mathbb{C}} x \rangle$ **for** $x :: \langle - :: \text{complex-normed-vector} \rangle$
 $\langle \text{proof} \rangle$

lemma *csubspace-has-basis*:
assumes $\langle \text{csubspace } S \rangle$
shows $\langle \exists B. \text{cindependent } B \wedge \text{cspan } B = S \rangle$
 $\langle \text{proof} \rangle$

lemma *inj-scaleC*:
fixes $A :: \langle 'a :: \text{complex-vector set} \rangle$
assumes $\langle c \neq 0 \rangle$
shows $\langle \text{inj-on } (\text{scaleC } c) \ A \rangle$
 $\langle \text{proof} \rangle$

7.2 Antilinear maps and friends

locale *antilinear* = *additive* f **for** $f :: 'a :: \text{complex-vector} \Rightarrow 'b :: \text{complex-vector} +$
assumes *scaleC*: $f\ (\text{scaleC } r\ x) = \text{cnj } r *_{\mathbb{C}} f\ x$

sublocale *antilinear* $\subseteq$ *linear*
 $\langle \text{proof} \rangle$

lemma *antilinear-imp-scaleC*:
fixes $D :: \text{complex} \Rightarrow 'a :: \text{complex-vector}$
assumes *antilinear* D

obtains d **where** $D = (\lambda x. \text{cnj } x *_{\mathbb{C}} d)$
 $\langle \text{proof} \rangle$

corollary *complex-antilinearD*:
fixes $f :: \text{complex} \Rightarrow \text{complex}$
assumes *antilinear* f **obtains** c **where** $f = (\lambda x. c * \text{cnj } x)$
 $\langle \text{proof} \rangle$

lemma *antilinearI*:
assumes $\bigwedge x y. f (x + y) = f x + f y$
and $\bigwedge c x. f (c *_{\mathbb{C}} x) = \text{cnj } c *_{\mathbb{C}} f x$
shows *antilinear* f
 $\langle \text{proof} \rangle$

lemma *antilinear-o-antilinear*: *antilinear* $f \Rightarrow \text{antilinear } g \Rightarrow \text{clinear } (g \circ f)$
 $\langle \text{proof} \rangle$

lemma *clinear-o-antilinear*: *antilinear* $f \Rightarrow \text{clinear } g \Rightarrow \text{antilinear } (g \circ f)$
 $\langle \text{proof} \rangle$

lemma *antilinear-o-clinear*: *clinear* $f \Rightarrow \text{antilinear } g \Rightarrow \text{antilinear } (g \circ f)$
 $\langle \text{proof} \rangle$

locale *bounded-antilinear* = *antilinear* f **for** $f :: 'a::\text{complex-normed-vector} \Rightarrow 'b::\text{complex-normed-vector} +$
assumes *bounded*: $\exists K. \forall x. \text{norm } (f x) \leq \text{norm } x * K$

lemma *bounded-antilinearI*:
assumes $\langle \bigwedge b1 b2. f (b1 + b2) = f b1 + f b2 \rangle$
assumes $\langle \bigwedge r b. f (r *_{\mathbb{C}} b) = \text{cnj } r *_{\mathbb{C}} f b \rangle$
assumes $\langle \forall x. \text{norm } (f x) \leq \text{norm } x * K \rangle$
shows *bounded-antilinear* f
 $\langle \text{proof} \rangle$

sublocale *bounded-antilinear* $\subseteq$ *real*: *bounded-linear*
— Gives access to all lemmas from *Real-Vector-Spaces.linear* using prefix *real*.
 $\langle \text{proof} \rangle$

lemma (**in** *bounded-antilinear*) *bounded-linear*: *bounded-linear* f
 $\langle \text{proof} \rangle$

lemma (**in** *bounded-antilinear*) *antilinear*: *antilinear* f
 $\langle \text{proof} \rangle$

lemma *bounded-antilinear-intro*:
assumes $\bigwedge x y. f (x + y) = f x + f y$
and $\bigwedge r x. f (\text{scaleC } r x) = \text{scaleC } (\text{cnj } r) (f x)$
and $\bigwedge x. \text{norm } (f x) \leq \text{norm } x * K$
shows *bounded-antilinear* f

$\langle \text{proof} \rangle$

lemma *bounded-antilinear-0*[simp]: $\langle \text{bounded-antilinear } (\lambda x. 0) \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-antilinear-0'*[simp]: $\langle \text{bounded-antilinear } 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *cnj-bounded-antilinear*[simp]: *bounded-antilinear cnj*
 $\langle \text{proof} \rangle$

lemma *bounded-antilinear-o-bounded-antilinear*:
 assumes *bounded-antilinear f*
 and *bounded-antilinear g*
 shows *bounded-clinear* $(\lambda x. f (g x))$
 $\langle \text{proof} \rangle$

lemma *bounded-antilinear-o-bounded-antilinear'*:
 assumes *bounded-antilinear f*
 and *bounded-antilinear g*
 shows *bounded-clinear* $(g \circ f)$
 $\langle \text{proof} \rangle$

lemma *bounded-antilinear-o-bounded-clinear*:
 assumes *bounded-antilinear f*
 and *bounded-clinear g*
 shows *bounded-antilinear* $(\lambda x. f (g x))$
 $\langle \text{proof} \rangle$

lemma *bounded-antilinear-o-bounded-clinear'*:
 assumes *bounded-clinear f*
 and *bounded-antilinear g*
 shows *bounded-antilinear* $(g \circ f)$
 $\langle \text{proof} \rangle$

lemma *bounded-clinear-o-bounded-antilinear*:
 assumes *bounded-clinear f*
 and *bounded-antilinear g*
 shows *bounded-antilinear* $(\lambda x. f (g x))$
 $\langle \text{proof} \rangle$

lemma *bounded-clinear-o-bounded-antilinear'*:
 assumes *bounded-antilinear f*
 and *bounded-clinear g*
 shows *bounded-antilinear* $(g \circ f)$
 $\langle \text{proof} \rangle$

lemma *bij-clinear-imp-inv-clinear*: *clinear* $(\text{inv } f)$
 if *a1*: *clinear f* **and** *a2*: *bij f*

$\langle proof \rangle$

locale *bounded-sesquilinear* =

fixes

prod :: 'a::complex-normed-vector $\Rightarrow$ 'b::complex-normed-vector $\Rightarrow$ 'c::complex-normed-vector
 (infixl <***> 70)

assumes *add-left*: *prod* (a + a') b = *prod* a b + *prod* a' b

and *add-right*: *prod* a (b + b') = *prod* a b + *prod* a b'

and *scaleC-left*: *prod* (r *_C a) b = (cnj r) *_C (*prod* a b)

and *scaleC-right*: *prod* a (r *_C b) = r *_C (*prod* a b)

and *bounded*: $\exists K. \forall a b. \text{norm } (\text{prod } a b) \leq \text{norm } a * \text{norm } b * K$

sublocale *bounded-sesquilinear* $\subseteq$ *real*: *bounded-bilinear*

— Gives access to all lemmas from *Real-Vector-Spaces.linear* using prefix *real*.

$\langle proof \rangle$

lemma (in *bounded-sesquilinear*) *bounded-bilinear*[simp]: *bounded-bilinear prod*

$\langle proof \rangle$

lemma (in *bounded-sesquilinear*) *bounded-antilinear-left*: *bounded-antilinear* ($\lambda a. \text{prod } a b$)

$\langle proof \rangle$

lemma (in *bounded-sesquilinear*) *bounded-clinear-right*: *bounded-clinear* ($\lambda b. \text{prod } a b$)

$\langle proof \rangle$

lemma (in *bounded-sesquilinear*) *comp1*:

assumes $\langle \text{bounded-clinear } g \rangle$

shows $\langle \text{bounded-sesquilinear } (\lambda x. \text{prod } (g x)) \rangle$

$\langle proof \rangle$

lemma (in *bounded-sesquilinear*) *comp2*:

assumes $\langle \text{bounded-clinear } g \rangle$

shows $\langle \text{bounded-sesquilinear } (\lambda x y. \text{prod } x (g y)) \rangle$

$\langle proof \rangle$

lemma (in *bounded-sesquilinear*) *comp*: *bounded-clinear* f $\implies$ *bounded-clinear* g $\implies$ *bounded-sesquilinear* ($\lambda x y. \text{prod } (f x) (g y)$)

$\langle proof \rangle$

lemma *bounded-clinear-const-scaleR*:

fixes c :: real

assumes $\langle \text{bounded-clinear } f \rangle$

shows $\langle \text{bounded-clinear } (\lambda x. c *_{\mathbb{R}} f x) \rangle$

$\langle proof \rangle$

lemma *bounded-linear-bounded-clinear*:

$\langle \text{bounded-linear } A \implies \forall c x. A (c *_C x) = c *_C A x \implies \text{bounded-clinear } A \rangle$
 $\langle \text{proof} \rangle$

lemma *comp-bounded-clinear*:

fixes $A :: \langle 'b :: \text{complex-normed-vector} \Rightarrow 'c :: \text{complex-normed-vector} \rangle$
and $B :: \langle 'a :: \text{complex-normed-vector} \Rightarrow 'b \rangle$
assumes $\langle \text{bounded-clinear } A \rangle$ **and** $\langle \text{bounded-clinear } B \rangle$
shows $\langle \text{bounded-clinear } (A \circ B) \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-sesquilinear-add*:

$\langle \text{bounded-sesquilinear } (\lambda x y. A x y + B x y) \rangle$ **if** $\langle \text{bounded-sesquilinear } A \rangle$ **and**
 $\langle \text{bounded-sesquilinear } B \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-sesquilinear-uminus*:

$\langle \text{bounded-sesquilinear } (\lambda x y. - A x y) \rangle$ **if** $\langle \text{bounded-sesquilinear } A \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-sesquilinear-diff*:

$\langle \text{bounded-sesquilinear } (\lambda x y. A x y - B x y) \rangle$ **if** $\langle \text{bounded-sesquilinear } A \rangle$ **and**
 $\langle \text{bounded-sesquilinear } B \rangle$
 $\langle \text{proof} \rangle$

lemmas *isCont-scaleC* [simp] =

bounded-bilinear.isCont [OF *bounded-cbilinear-scaleC* [THEN *bounded-cbilinear.bounded-bilinear*]]

lemma *bounded-sesquilinear-0* [simp]: $\langle \text{bounded-sesquilinear } (\lambda -. 0) \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-sesquilinear-0'* [simp]: $\langle \text{bounded-sesquilinear } 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-sesquilinear-apply-bounded-clinear*: $\langle \text{bounded-clinear } (f x) \rangle$ **if** $\langle \text{bounded-sesquilinear } f \rangle$
 $\langle \text{proof} \rangle$

lemma *antilinear-diff*:

assumes $\langle \text{antilinear } f \rangle$ **and** $\langle \text{antilinear } g \rangle$
shows $\langle \text{antilinear } (\lambda x. f x - g x) \rangle$
 $\langle \text{proof} \rangle$

7.3 Misc 2

lemma *summable-on-scaleC-left* [intro]:

fixes $c :: \langle 'a :: \text{complex-normed-vector} \rangle$
assumes $c \neq 0 \implies f \text{ summable-on } A$
shows $(\lambda x. f x *_C c) \text{ summable-on } A$

$\langle \text{proof} \rangle$

lemma *summable-on-scaleC-right* [intro]:
fixes $f :: \langle 'a \Rightarrow 'b :: \text{complex-normed-vector} \rangle$
assumes $c \neq 0 \implies f \text{ summable-on } A$
shows $(\lambda x. c *_C f x) \text{ summable-on } A$
 $\langle \text{proof} \rangle$

lemma *infsum-scaleC-left*:
fixes $c :: \langle 'a :: \text{complex-normed-vector} \rangle$
assumes $c \neq 0 \implies f \text{ summable-on } A$
shows $\text{infsum } (\lambda x. f x *_C c) A = \text{infsum } f A *_C c$
 $\langle \text{proof} \rangle$

lemma *infsum-scaleC-right*:
fixes $f :: \langle 'a \Rightarrow 'b :: \text{complex-normed-vector} \rangle$
shows $\text{infsum } (\lambda x. c *_C f x) A = c *_C \text{infsum } f A$
 $\langle \text{proof} \rangle$

lemma *has-sum-scaleC-right*:
fixes $f :: \langle 'a \Rightarrow 'b :: \text{complex-normed-vector} \rangle$
assumes $\langle f \text{ has-sum } s \rangle A$
shows $\langle (\lambda x. c *_C f x) \text{ has-sum } c *_C s \rangle A$
 $\langle \text{proof} \rangle$

lemmas *sums-of-complex* = *bounded-linear.sums* [OF *bounded-clinear-of-complex* [THEN *bounded-clinear.bounded-linear*]]
lemmas *summable-of-complex* = *bounded-linear.summable* [OF *bounded-clinear-of-complex* [THEN *bounded-clinear.bounded-linear*]]
lemmas *suminf-of-complex* = *bounded-linear.suminf* [OF *bounded-clinear-of-complex* [THEN *bounded-clinear.bounded-linear*]]

lemmas *sums-scaleC-left* = *bounded-linear.sums* [OF *bounded-clinear-scaleC-left* [THEN *bounded-clinear.bounded-linear*]]
lemmas *summable-scaleC-left* = *bounded-linear.summable* [OF *bounded-clinear-scaleC-left* [THEN *bounded-clinear.bounded-linear*]]
lemmas *suminf-scaleC-left* = *bounded-linear.suminf* [OF *bounded-clinear-scaleC-left* [THEN *bounded-clinear.bounded-linear*]]

lemmas *sums-scaleC-right* = *bounded-linear.sums* [OF *bounded-clinear-scaleC-right* [THEN *bounded-clinear.bounded-linear*]]
lemmas *summable-scaleC-right* = *bounded-linear.summable* [OF *bounded-clinear-scaleC-right* [THEN *bounded-clinear.bounded-linear*]]
lemmas *suminf-scaleC-right* = *bounded-linear.suminf* [OF *bounded-clinear-scaleC-right* [THEN *bounded-clinear.bounded-linear*]]

lemma *closed-scaleC*:
fixes $S :: \langle 'a :: \text{complex-normed-vector set} \rangle$ **and** $a :: \text{complex}$

```

assumes  $\langle \text{closed } S \rangle$ 
shows  $\langle \text{closed } ((*_C) a \text{ ' } S) \rangle$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma closure-scaleC:
  fixes  $S :: \langle 'a :: \text{complex-normed-vector set} \rangle$ 
  shows  $\langle \text{closure } ((*_C) a \text{ ' } S) = (*_C) a \text{ ' closure } S \rangle$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma onorm-scalarC:
  fixes  $f :: \langle 'a :: \text{complex-normed-vector} \Rightarrow 'b :: \text{complex-normed-vector} \rangle$ 
  assumes  $a1 :: \langle \text{bounded-clinear } f \rangle$ 
  shows  $\langle \text{onorm } (\lambda x. r *_C (f x)) = (cmod r) * \text{onorm } f \rangle$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma onorm-scaleC-left-lemma:
  fixes  $f :: 'a :: \text{complex-normed-vector}$ 
  assumes  $r :: \text{bounded-clinear } r$ 
  shows  $\text{onorm } (\lambda x. r x *_C f) \leq \text{onorm } r * \text{norm } f$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma onorm-scaleC-left:
  fixes  $f :: 'a :: \text{complex-normed-vector}$ 
  assumes  $f :: \text{bounded-clinear } r$ 
  shows  $\text{onorm } (\lambda x. r x *_C f) = \text{onorm } r * \text{norm } f$ 
 $\langle \text{proof} \rangle$ 

```

```

lemma compact-scaleC:
  fixes  $s :: 'a :: \text{complex-normed-vector set}$ 
  assumes compact s
  shows compact (scaleC c ' s)
 $\langle \text{proof} \rangle$ 

```

7.4 Finite dimension and canonical basis

```

lemma vector-finitely-spanned:
  assumes  $\langle z \in \text{cspan } T \rangle$ 
  shows  $\langle \exists S. \text{finite } S \wedge S \subseteq T \wedge z \in \text{cspan } S \rangle$ 
 $\langle \text{proof} \rangle$ 

```

$\langle ML \rangle$

```

class cfinite-dim = complex-vector +
  assumes cfinutely-spanned:  $\exists S :: 'a \text{ set. finite } S \wedge \text{cspan } S = UNIV$ 

```

```

class basis-enum = complex-vector +
  fixes canonical-basis ::  $\langle 'a \text{ list} \rangle$ 
  and canonical-basis-length ::  $\langle 'a \text{ itself} \Rightarrow \text{nat} \rangle$ 
  assumes distinct-canonical-basis[simp]:

```

```

distinct canonical-basis
and is-cindependent-set[simp]:
  cindependent (set canonical-basis)
and is-generator-set[simp]:
  cspan (set canonical-basis) = UNIV
and canonical-basis-length:
   $\langle \text{canonical-basis-length } \text{TYPE}('a) = \text{length } \text{canonical-basis} \rangle$ 

```

$\langle ML \rangle$

```

instantiation complex :: basis-enum begin
definition canonical-basis = [1::complex]
definition  $\langle \text{canonical-basis-length } (-::\text{complex itself}) = 1 \rangle$ 
instance
 $\langle \text{proof} \rangle$ 
end

lemma cdim-UNIV-basis-enum[simp]:  $\langle \text{cdim } (\text{UNIV}::'a::\text{basis-enum set}) = \text{length } (\text{canonical-basis}::'a \text{ list}) \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma finite-basis:  $\exists \text{basis}::'a::\text{cfinite-dim set. finite basis} \wedge \text{cindependent basis} \wedge \text{cspan basis} = \text{UNIV}$ 
 $\langle \text{proof} \rangle$ 

instance basis-enum  $\subseteq$  cfinite-dim
 $\langle \text{proof} \rangle$ 

lemma cindependent-cfinite-dim-finite:
  assumes  $\langle \text{cindependent } (S::'a::\text{cfinite-dim set}) \rangle$ 
  shows  $\langle \text{finite } S \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma cfinite-dim-finite-subspace-basis:
  assumes  $\langle \text{csubspace } X \rangle$ 
  shows  $\exists \text{basis}::'a::\text{cfinite-dim set. finite basis} \wedge \text{cindependent basis} \wedge \text{cspan basis} = X$ 
 $\langle \text{proof} \rangle$ 

```

The following auxiliary lemma (*finite-span-complete-aux*) shows more or less the same as *finite-span-representation-bounded*, *finite-span-complete* below (see there for an intuition about the mathematical content of the lemmas). However, there is one difference: Here we additionally assume here that there is a bijection *rep/abs* between a finite type *'basis* and the set *B*. This is needed to be able to use results about euclidean spaces that are formulated w.r.t. the type class *finite*

Since we anyway assume that *B* is finite, this added assumption does not

make the lemma weaker. However, we cannot derive the existence of *'basis* inside the proof (HOL does not support such reasoning). Therefore we have the type *'basis* as an explicit assumption and remove it using *internalize-sort* after the proof.

```

lemma finite-span-complete-aux:
  fixes b :: 'b::real-normed-vector and B :: 'b set
    and rep :: 'basis::finite  $\Rightarrow$  'b and abs :: 'b  $\Rightarrow$  'basis
  assumes t: type-definition rep abs B
    and t1: finite B and t2: b ∈ B and t3: independent B
  shows  $\exists D > 0. \forall \psi. \text{norm} (\text{representation } B \ \psi \ b) \leq \text{norm } \psi * D$ 
    and complete (span B)
  <proof>

```

```

lemma finite-span-complete[simp]:
  fixes A :: 'a::real-normed-vector set
  assumes finite A
  shows complete (span A)

```

The span of a finite set is complete.

<proof>

```

lemma finite-span-representation-bounded:
  fixes B :: 'a::real-normed-vector set
  assumes finite B and independent B
  shows  $\exists D > 0. \forall \psi \ b. \text{abs} (\text{representation } B \ \psi \ b) \leq \text{norm } \psi * D$ 

```

Assume *B* is a finite linear independent set of vectors (in a real normed vector space). Let α_b^ψ be the coefficients of ψ expressed as a linear combination over *B*. Then α is uniformly cblinfun (i.e., $|\alpha_b^\psi| \leq D \|\psi\|$ for some *D* independent of ψ, b).

(This also holds when *b* is not in the span of *B* because of the way *real-vector.representation* is defined in this corner case.)

<proof>

hide-fact *finite-span-complete-aux*

```

lemma finite-cspan-complete[simp]:
  fixes B :: 'a::complex-normed-vector set
  assumes finite B
  shows complete (cspan B)
  <proof>

```

```

lemma finite-span-closed[simp]:
  fixes B :: 'a::real-normed-vector set
  assumes finite B
  shows closed (real-vector.span B)

```

⟨proof⟩

lemma *finite-cspan-closed*[simp]:
fixes $S :: \langle 'a :: \text{complex-normed-vector set} \rangle$
assumes $a1 :: \langle \text{finite } S \rangle$
shows $\langle \text{closed } (\text{cspan } S) \rangle$
 ⟨proof⟩

lemma *closure-finite-cspan*:
fixes $T :: \langle 'a :: \text{complex-normed-vector set} \rangle$
assumes $\langle \text{finite } T \rangle$
shows $\langle \text{closure } (\text{cspan } T) = \text{cspan } T \rangle$
 ⟨proof⟩

lemma *finite-cspan-crepresentation-bounded*:
fixes $B :: 'a :: \text{complex-normed-vector set}$
assumes $a1 :: \text{finite } B$ **and** $a2 :: \text{cindependent } B$
shows $\exists D > 0. \forall \psi \ b. \text{cmod } (\text{crepresentation } B \ \psi \ b) \leq \text{norm } \psi * D$
 ⟨proof⟩

lemma *bounded-clinear-finite-dim*[simp]:
fixes $f :: \langle 'a :: \{ \text{cfinite-dim, complex-normed-vector} \} \Rightarrow 'b :: \text{complex-normed-vector} \rangle$
assumes $\langle \text{clinear } f \rangle$
shows $\langle \text{bounded-clinear } f \rangle$
 ⟨proof⟩
include *norm-syntax*
 ⟨proof⟩

lemma *summable-on-scaleR-left-converse*:
 — This result has nothing to do with the bounded operator library but it uses *finite-span-closed* so it is proven here.
fixes $f :: \langle 'b \Rightarrow \text{real} \rangle$
and $c :: \langle 'a :: \text{real-normed-vector} \rangle$
assumes $\langle c \neq 0 \rangle$
assumes $\langle (\lambda x. f \ x *_{\mathbb{R}} c) \text{ summable-on } A \rangle$
shows $\langle f \text{ summable-on } A \rangle$
 ⟨proof⟩

lemma *infsum-scaleR-left*:
 — This result has nothing to do with the bounded operator library but it uses *finite-span-closed* so it is proven here.
 It is a strengthening of *infsum-scaleR-left*.
fixes $c :: \langle 'a :: \text{real-normed-vector} \rangle$
shows $\text{infsum } (\lambda x. f \ x *_{\mathbb{R}} c) \ A = \text{infsum } f \ A *_{\mathbb{R}} c$
 ⟨proof⟩

lemma *infsum-of-real*:

shows $\langle (\sum_{\infty x \in A. \text{ of-real } (f \ x)} :: 'b::\{\text{real-normed-vector}, \text{real-algebra-1}\}) = \text{of-real } (\sum_{\infty x \in A. f \ x}) \rangle$
 — This result has nothing to do with the bounded operator library but it uses *finite-span-closed* so it is proven here.
 $\langle \text{proof} \rangle$

definition $\langle \text{cfinite-dim } S \longleftrightarrow (\exists B. \text{finite } B \wedge S \subseteq \text{cspan } B) \rangle$

lemma *cspan-finite-dim*[intro]: $\langle \text{cfinite-dim } (\text{cspan } B) \rangle$ **if** $\langle \text{finite } B \rangle$
 $\langle \text{proof} \rangle$

lemma *cfinite-dim-subspace-has-basis*:
assumes $\langle \text{cfinite-dim } S \rangle$ **and** $\langle \text{csubspace } S \rangle$
shows $\langle \exists B. \text{finite } B \wedge \text{cindependent } B \wedge \text{cspan } B = S \rangle$
 $\langle \text{proof} \rangle$

7.5 Closed subspaces

lemma *csubspace-INF*[simp]: $(\bigwedge x. x \in A \implies \text{csubspace } x) \implies \text{csubspace } (\bigcap A)$
 $\langle \text{proof} \rangle$

locale *closed-csubspace* =
fixes $A::('a::\{\text{complex-vector}, \text{topological-space}\}) \text{ set}$
assumes *subspace*: $\text{csubspace } A$
assumes *closed*: $\text{closed } A$

declare *closed-csubspace.subspace*[simp]

lemma *closure-is-csubspace*[simp]:
fixes $A::('a::\{\text{complex-normed-vector}\}) \text{ set}$
assumes $\langle \text{csubspace } A \rangle$
shows $\langle \text{csubspace } (\text{closure } A) \rangle$
 $\langle \text{proof} \rangle$

lemma *csubspace-set-plus*:
assumes $\langle \text{csubspace } A \rangle$ **and** $\langle \text{csubspace } B \rangle$
shows $\langle \text{csubspace } (A + B) \rangle$
 $\langle \text{proof} \rangle$

lemma *closed-csubspace-0*[simp]:
 $\text{closed-csubspace } (\{0\} :: ('a::\{\text{complex-vector}, \text{t1-space}\}) \text{ set})$
 $\langle \text{proof} \rangle$

lemma *closed-csubspace-UNIV*[simp]: $\text{closed-csubspace } (\text{UNIV}::('a::\{\text{complex-vector}, \text{topological-space}\}) \text{ set})$
 $\langle \text{proof} \rangle$

lemma *closed-csubspace-inter*[simp]:

```

    assumes closed-csubspace A and closed-csubspace B
    shows closed-csubspace (A ∩ B)
  <proof>

lemma closed-csubspace-INF[simp]:
  assumes a1:  $\forall A \in \mathcal{A}. \text{closed-csubspace } A$ 
  shows closed-csubspace ( $\bigcap \mathcal{A}$ )
  <proof>

typedef (overloaded) ('a::complex-vector, topological-space)
  csubspace = ⟨{S::'a set. closed-csubspace S}⟩
  morphisms space-as-set Abs-ccsubspace
  <proof>

setup-lifting type-definition-ccsubspace

lemma csubspace-space-as-set[simp]: csubspace (space-as-set S)
  <proof>

lemma closed-space-as-set[simp]: closed (space-as-set S)
  <proof>

lemma zero-space-as-set[simp]:  $0 \in \text{space-as-set } A$ 
  <proof>

instantiation ccsubspace :: (complex-normed-vector) scaleC begin
lift-definition scaleC-ccsubspace :: complex  $\Rightarrow$  'a ccsubspace  $\Rightarrow$  'a ccsubspace is
   $\lambda c S. (*_C) c \text{ ' } S$ 
  <proof>

lift-definition scaleR-ccsubspace :: real  $\Rightarrow$  'a ccsubspace  $\Rightarrow$  'a ccsubspace is
   $\lambda c S. (*_R) c \text{ ' } S$ 
  <proof>

instance
  <proof>
end

instantiation ccsubspace :: ({complex-vector, t1-space}) bot begin
lift-definition bot-ccsubspace :: 'a ccsubspace is ⟨{0}⟩
  <proof>
instance<proof>
end

lemma zero-cblinfun-image[simp]:  $0 *_C S = \text{bot}$  for S :: - ccsubspace
  <proof>

lemma csubspace-scaleC-invariant:
```

fixes $a\ S$
assumes $\langle a \neq 0 \rangle$ **and** $\langle \text{ccsubspace } S \rangle$
shows $\langle (*_C) a \cdot S = S \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{ccsubspace-scaleC-invariant}[simp]$: $a \neq 0 \implies a *_C S = S$ **for** $S :: - \text{ccsubspace}$
 $\langle \text{proof} \rangle$

instantiation $\text{ccsubspace} :: (\{\text{complex-vector}, \text{topological-space}\}) \text{ top}$
begin
lift-definition $\text{top-ccsubspace} :: \langle 'a \text{ ccsubspace} \rangle$ **is** $\langle \text{UNIV} \rangle$
 $\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$
end

lemma $\text{space-as-set-bot}[simp]$: $\langle \text{space-as-set bot} = \{0\} \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{ccsubspace-top-not-bot}[simp]$:
 $(\text{top} :: 'a :: \{\text{complex-vector}, \text{t1-space}, \text{not-singleton}\} \text{ ccsubspace}) \neq \text{bot}$
 $\langle \text{proof} \rangle$

lemma $\text{ccsubspace-bot-not-top}[simp]$:
 $(\text{bot} :: 'a :: \{\text{complex-vector}, \text{t1-space}, \text{not-singleton}\} \text{ ccsubspace}) \neq \text{top}$
 $\langle \text{proof} \rangle$

instantiation $\text{ccsubspace} :: (\{\text{complex-vector}, \text{topological-space}\}) \text{ Inf}$
begin
lift-definition $\text{Inf-ccsubspace} :: \langle 'a \text{ ccsubspace set} \Rightarrow 'a \text{ ccsubspace} \rangle$
is $\langle \lambda S. \bigcap S \rangle$
 $\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$
end

lift-definition $\text{ccspan} :: 'a :: \text{complex-normed-vector set} \Rightarrow 'a \text{ ccsubspace}$
is $\lambda G. \text{closure } (\text{cspan } G)$
 $\langle \text{proof} \rangle$

lemma ccspan-superset :
 $\langle A \subseteq \text{space-as-set } (\text{ccspan } A) \rangle$
for $A :: \langle 'a :: \text{complex-normed-vector set} \rangle$
 $\langle \text{proof} \rangle$

lemma *ccspan-superset'*: $\langle x \in X \implies x \in \text{space-as-set } (\text{ccspan } X) \rangle$
 $\langle \text{proof} \rangle$

lemma *ccspan-canonical-basis[simp]*: $\text{ccspan } (\text{set canonical-basis}) = \text{top}$
 $\langle \text{proof} \rangle$

lemma *ccspan-Inf-def*: $\langle \text{ccspan } A = \text{Inf } \{S. A \subseteq \text{space-as-set } S\} \rangle$
for $A :: \langle 'a :: \text{cbanach} \rangle \text{ set} \rangle$
 $\langle \text{proof} \rangle$

lemma *cspan-singleton-scaleC[simp]*: $(a :: \text{complex}) \neq 0 \implies \text{cspan } \{ a *_C \psi \} =$
 $\text{cspan } \{ \psi \}$
for $\psi :: 'a :: \text{complex-vector}$
 $\langle \text{proof} \rangle$

lemma *closure-is-closed-csubspace[simp]*:
fixes $S :: \langle 'a :: \text{complex-normed-vector set} \rangle$
assumes $\langle \text{csubspace } S \rangle$
shows $\langle \text{closed-csubspace } (\text{closure } S) \rangle$
 $\langle \text{proof} \rangle$

lemma *ccspan-singleton-scaleC[simp]*: $(a :: \text{complex}) \neq 0 \implies \text{ccspan } \{ a *_C \psi \} =$
 $\text{ccspan } \{ \psi \}$
 $\langle \text{proof} \rangle$

lemma *clinear-continuous-at*:
assumes $\langle \text{bounded-clinear } f \rangle$
shows $\langle \text{isCont } f \ x \rangle$
 $\langle \text{proof} \rangle$

lemma *clinear-continuous-within*:
assumes $\langle \text{bounded-clinear } f \rangle$
shows $\langle \text{continuous } (\text{at } x \text{ within } s) \ f \rangle$
 $\langle \text{proof} \rangle$

lemma *antilinear-continuous-at*:
assumes $\langle \text{bounded-antilinear } f \rangle$
shows $\langle \text{isCont } f \ x \rangle$
 $\langle \text{proof} \rangle$

lemma *antilinear-continuous-within*:
assumes $\langle \text{bounded-antilinear } f \rangle$
shows $\langle \text{continuous } (\text{at } x \text{ within } s) \ f \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-clinear-eq-on-closure*:
fixes $A \ B :: 'a :: \text{complex-normed-vector} \Rightarrow 'b :: \text{complex-normed-vector}$
assumes $\langle \text{bounded-clinear } A \rangle$ **and** $\langle \text{bounded-clinear } B \rangle$ **and**
 $\text{eq: } \langle \bigwedge x. x \in G \implies A \ x = B \ x \rangle$ **and** $t: \langle t \in \text{closure } (\text{cspan } G) \rangle$

shows $\langle A \ t = B \ t \rangle$
 $\langle proof \rangle$

instantiation $ccsubspace :: (\{complex-vector, topological-space\}) \text{ order}$
begin
lift-definition $less-eq-ccsubspace :: \langle 'a \ ccsubspace \Rightarrow 'a \ ccsubspace \Rightarrow bool \rangle$
is $\langle (\subseteq) \rangle \langle proof \rangle$
declare $less-eq-ccsubspace-def[code \ del]$
lift-definition $less-ccsubspace :: \langle 'a \ ccsubspace \Rightarrow 'a \ ccsubspace \Rightarrow bool \rangle$
is $\langle (\subset) \rangle \langle proof \rangle$
declare $less-ccsubspace-def[code \ del]$
instance
 $\langle proof \rangle$
end

lemma $ccspan-leqI$:
assumes $\langle M \subseteq space-as-set \ S \rangle$
shows $\langle ccspan \ M \leq S \rangle$
 $\langle proof \rangle$

lemma $ccspan-mono$:
assumes $\langle A \subseteq B \rangle$
shows $\langle ccspan \ A \leq ccspan \ B \rangle$
 $\langle proof \rangle$

lemma $ccsubspace-leI$:
assumes $t1: space-as-set \ A \subseteq space-as-set \ B$
shows $A \leq B$
 $\langle proof \rangle$

lemma $ccspan-of-empty[simp]$: $ccspan \ \{\} = bot$
 $\langle proof \rangle$

instantiation $ccsubspace :: (\{complex-vector, topological-space\}) \text{ inf}$ **begin**
lift-definition $inf-ccsubspace :: 'a \ ccsubspace \Rightarrow 'a \ ccsubspace \Rightarrow 'a \ ccsubspace$
is $(\cap) \langle proof \rangle$
instance $\langle proof \rangle$ **end**

lemma $space-as-set-inf[simp]$: $space-as-set \ (A \sqcap B) = space-as-set \ A \cap space-as-set \ B$
 $\langle proof \rangle$

instantiation $ccsubspace :: (\{complex-vector, topological-space\}) \text{ order-top}$ **begin**
instance
 $\langle proof \rangle$
end

instantiation *ccsubspace* :: (*{complex-vector,t1-space}*) *order-bot* **begin**
instance
 $\langle proof \rangle$
end

instantiation *ccsubspace* :: (*{complex-vector,topological-space}*) *semilattice-inf* **begin**
instance
 $\langle proof \rangle$
end

instantiation *ccsubspace* :: (*{complex-vector,t1-space}*) *zero* **begin**
definition *zero-ccsubspace* :: 'a *ccsubspace* **where** [*simp*]: *zero-ccsubspace* = *bot*
lemma *zero-ccsubspace-transfer*[*transfer-rule*]: $\langle pcr\text{-}ccsubspace\ (=)\ \{0\}\ 0 \rangle$
 $\langle proof \rangle$
instance $\langle proof \rangle$
end

lemma *ccspan-0*[*simp*]: $\langle ccspan\ \{0\} = 0 \rangle$
 $\langle proof \rangle$

definition $\langle rel\text{-}ccsubspace\ R\ x\ y = rel\text{-}set\ R\ (space\text{-}as\text{-}set\ x)\ (space\text{-}as\text{-}set\ y) \rangle$

lemma *left-unique-rel-ccsubspace*[*transfer-rule*]: $\langle left\text{-}unique\ (rel\text{-}ccsubspace\ R) \rangle$ **if**
 $\langle left\text{-}unique\ R \rangle$
 $\langle proof \rangle$

lemma *right-unique-rel-ccsubspace*[*transfer-rule*]: $\langle right\text{-}unique\ (rel\text{-}ccsubspace\ R) \rangle$
if $\langle right\text{-}unique\ R \rangle$
 $\langle proof \rangle$

lemma *bi-unique-rel-ccsubspace*[*transfer-rule*]: $\langle bi\text{-}unique\ (rel\text{-}ccsubspace\ R) \rangle$ **if** $\langle bi\text{-}unique\ R \rangle$
 $\langle proof \rangle$

lemma *converse-rel-ccsubspace*: $\langle conversep\ (rel\text{-}ccsubspace\ R) = rel\text{-}ccsubspace\ (conversep\ R) \rangle$
 $\langle proof \rangle$

lemma *space-as-set-top*[*simp*]: $\langle space\text{-}as\text{-}set\ top = UNIV \rangle$
 $\langle proof \rangle$

lemma *ccsubspace-eqI*:
assumes $\langle \bigwedge x. x \in space\text{-}as\text{-}set\ S \longleftrightarrow x \in space\text{-}as\text{-}set\ T \rangle$
shows $\langle S = T \rangle$
 $\langle proof \rangle$

lemma *ccspan-remove-0*: $\langle \text{ccspan } (A - \{0\}) = \text{ccspan } A \rangle$
 $\langle \text{proof} \rangle$

lemma *sgn-in-spaceD*: $\langle \psi \in \text{space-as-set } A \rangle$ **if** $\langle \text{sgn } \psi \in \text{space-as-set } A \rangle$ **and** $\langle \psi \neq 0 \rangle$
for $\psi :: \langle - :: \text{complex-normed-vector} \rangle$
 $\langle \text{proof} \rangle$

lemma *sgn-in-spaceI*: $\langle \text{sgn } \psi \in \text{space-as-set } A \rangle$ **if** $\langle \psi \in \text{space-as-set } A \rangle$
for $\psi :: \langle - :: \text{complex-normed-vector} \rangle$
 $\langle \text{proof} \rangle$

lemma *ccsubspace-leI-unit*:
fixes $A B :: \langle - :: \text{complex-normed-vector } \text{ccsubspace} \rangle$
assumes $\bigwedge \psi. \text{norm } \psi = 1 \implies \psi \in \text{space-as-set } A \implies \psi \in \text{space-as-set } B$
shows $A \leq B$
 $\langle \text{proof} \rangle$

lemma *kernel-is-closed-csubspace[simp]*:
assumes $a1: \text{bounded-clinear } f$
shows $\text{closed-csubspace } (f - \{0\})$
 $\langle \text{proof} \rangle$

lemma *ccspan-closure[simp]*: $\langle \text{ccspan } (\text{closure } X) = \text{ccspan } X \rangle$
 $\langle \text{proof} \rangle$

lemma *ccspan-finite*: $\langle \text{space-as-set } (\text{ccspan } X) = \text{cspan } X \rangle$ **if** $\langle \text{finite } X \rangle$
 $\langle \text{proof} \rangle$

lemma *ccspan-UNIV[simp]*: $\langle \text{ccspan } \text{UNIV} = \top \rangle$
 $\langle \text{proof} \rangle$

lemma *infsum-in-closed-csubspaceI*:
assumes $\langle \bigwedge x. x \in X \implies f x \in A \rangle$
assumes $\langle \text{closed-csubspace } A \rangle$
shows $\langle \text{infsum } f X \in A \rangle$
 $\langle \text{proof} \rangle$

lemma *closed-csubspace-space-as-set[simp]*: $\langle \text{closed-csubspace } (\text{space-as-set } X) \rangle$
 $\langle \text{proof} \rangle$

lift-definition *finite-dim-ccsubspace* :: $\langle 'a :: \text{complex-normed-vector } \text{ccsubspace} \Rightarrow \text{bool} \rangle$ **is** *cfinite-dim* $\langle \text{proof} \rangle$

lemma *ccspan-finite-dim[intro]*: $\langle \text{finite-dim-ccsubspace } (\text{ccspan } B) \rangle$ **if** $\langle \text{finite } B \rangle$
 $\langle \text{proof} \rangle$

lemma *finite-dim-ccsubspace-zero[iff]*: $\langle \text{finite-dim-ccsubspace } 0 \rangle$

$\langle \text{proof} \rangle$

lemma *finite-dim-ccsubspace-bot*[iff]: $\langle \text{finite-dim-ccsubspace } \perp \rangle$
 $\langle \text{proof} \rangle$

lemma *ccsubspace-contains-unit*:
 assumes $\langle E \neq \perp \rangle$
 shows $\langle \exists h \in \text{space-as-set } E. \text{ norm } h = 1 \rangle$
 $\langle \text{proof} \rangle$

7.6 Closed sums

definition *closed-sum*:: $\langle 'a::\{\text{semigroup-add}, \text{topological-space}\} \text{ set} \Rightarrow 'a \text{ set} \Rightarrow 'a \text{ set} \rangle$ **where**
 $\langle \text{closed-sum } A \ B = \text{closure } (A + B) \rangle$

notation *closed-sum* (**infixl** $\langle +_M \rangle$ 65)

lemma *closed-sum-comm*: $\langle A +_M B = B +_M A \rangle$ **for** $A \ B :: \text{--ab-semigroup-add}$
 $\langle \text{proof} \rangle$

lemma *closed-sum-left-subset*: $\langle 0 \in B \implies A \subseteq A +_M B \rangle$ **for** $A \ B :: \text{--monoid-add}$
 $\langle \text{proof} \rangle$

lemma *closed-sum-right-subset*: $\langle 0 \in A \implies B \subseteq A +_M B \rangle$ **for** $A \ B :: \text{--monoid-add}$
 $\langle \text{proof} \rangle$

lemma *finite-cspan-closed-csubspace*:
 assumes *finite* ($S::'a::\text{complex-normed-vector set}$)
 shows *closed-csubspace* (*cspan* S)
 $\langle \text{proof} \rangle$

lemma *closed-sum-is-sup*:
 fixes $A \ B \ C::\langle 'a::\{\text{complex-vector}, \text{topological-space}\} \rangle \text{ set}$
 assumes $\langle \text{closed-csubspace } C \rangle$
 assumes $\langle A \subseteq C \rangle$ **and** $\langle B \subseteq C \rangle$
 shows $\langle (A +_M B) \subseteq C \rangle$
 $\langle \text{proof} \rangle$

lemma *closed-subspace-closed-sum*:
 fixes $A \ B::\langle 'a::\text{complex-normed-vector} \rangle \text{ set}$
 assumes $a1::\langle \text{csubspace } A \rangle$ **and** $a2::\langle \text{csubspace } B \rangle$
 shows $\langle \text{closed-csubspace } (A +_M B) \rangle$
 $\langle \text{proof} \rangle$

lemma *closed-sum-assoc*:
 fixes $A \ B \ C::'a::\text{real-normed-vector set}$

shows $\langle A +_M (B +_M C) = (A +_M B) +_M C \rangle$
 $\langle proof \rangle$

lemma *closed-sum-zero-left*[simp]:
fixes $A :: \langle 'a :: \{monoid-add, topological-space\} \rangle set$
shows $\langle \{0\} +_M A = closure\ A \rangle$
 $\langle proof \rangle$

lemma *closed-sum-zero-right*[simp]:
fixes $A :: \langle 'a :: \{monoid-add, topological-space\} \rangle set$
shows $\langle A +_M \{0\} = closure\ A \rangle$
 $\langle proof \rangle$

lemma *closed-sum-closure-right*[simp]:
fixes $A\ B :: \langle 'a :: real-normed-vector\ set \rangle$
shows $\langle A +_M closure\ B = A +_M B \rangle$
 $\langle proof \rangle$

lemma *closed-sum-closure-left*[simp]:
fixes $A\ B :: \langle 'a :: real-normed-vector\ set \rangle$
shows $\langle closure\ A +_M B = A +_M B \rangle$
 $\langle proof \rangle$

lemma *closed-sum-mono-left*:
assumes $\langle A \subseteq B \rangle$
shows $\langle A +_M C \subseteq B +_M C \rangle$
 $\langle proof \rangle$

lemma *closed-sum-mono-right*:
assumes $\langle A \subseteq B \rangle$
shows $\langle C +_M A \subseteq C +_M B \rangle$
 $\langle proof \rangle$

instantiation *ccsubspace* :: (complex-normed-vector) sup **begin**

lift-definition *sup-ccsubspace* :: $'a\ ccsubspace \Rightarrow 'a\ ccsubspace \Rightarrow 'a\ ccsubspace$

— Note that $A + B$ would not be a closed subspace, we need the closure. See, e.g., <https://math.stackexchange.com/a/1786792/403528>.

is $\lambda A\ B :: 'a\ set. A +_M B$
 $\langle proof \rangle$

instance $\langle proof \rangle$

end

lemma *closed-sum-cspan*[simp]:
shows $\langle cspan\ X +_M cspan\ Y = closure\ (cspan\ (X \cup Y)) \rangle$
 $\langle proof \rangle$

lemma *closure-image-closed-sum*:
assumes $\langle bounded-linear\ U \rangle$

shows $\langle \text{closure } (U \text{ ' } (A +_M B)) = \text{closure } (U \text{ ' } A) +_M \text{closure } (U \text{ ' } B) \rangle$
 $\langle \text{proof} \rangle$

lemma *ccspan-union*: $\text{ccspan } A \sqcup \text{ccspan } B = \text{ccspan } (A \cup B)$
 $\langle \text{proof} \rangle$

instantiation *ccsubspace* :: (*complex-normed-vector*) *Sup*
begin
lift-definition *Sup-ccsubspace*:: $\langle 'a \text{ ccsubspace set} \Rightarrow 'a \text{ ccsubspace} \rangle$
is $\langle \lambda S. \text{closure } (\text{complex-vector.span } (\text{Union } S)) \rangle$
 $\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$
end

instance *ccsubspace* :: ($\{\text{complex-normed-vector}\}$) *semilattice-sup*
 $\langle \text{proof} \rangle$

instance *ccsubspace* :: (*complex-normed-vector*) *complete-lattice*
 $\langle \text{proof} \rangle$

instantiation *ccsubspace* :: (*complex-normed-vector*) *comm-monoid-add* **begin**
definition *plus-ccsubspace* :: $'a \text{ ccsubspace} \Rightarrow - \Rightarrow -$
where [*simp*]: *plus-ccsubspace* = *sup*
instance
 $\langle \text{proof} \rangle$
end

lemma *SUP-ccspan*: $\langle (\text{SUP } x \in X. \text{ccspan } (S \ x)) = \text{ccspan } (\bigcup x \in X. S \ x) \rangle$
 $\langle \text{proof} \rangle$

lemma *ccsubspace-plus-sup*: $y \leq x \implies z \leq x \implies y + z \leq x$
for $x \ y \ z :: 'a :: \text{complex-normed-vector ccsubspace}$
 $\langle \text{proof} \rangle$

lemma *ccsubspace-Sup-empty*: $\text{Sup } \{\} = (0 :: - \text{ ccsubspace})$
 $\langle \text{proof} \rangle$

lemma *ccsubspace-add-right-incr*[*simp*]: $a \leq a + c$ **for** $a :: - \text{ ccsubspace}$
 $\langle \text{proof} \rangle$

lemma *ccsubspace-add-left-incr*[*simp*]: $a \leq c + a$ **for** $a :: - \text{ ccsubspace}$
 $\langle \text{proof} \rangle$

lemma *sum-bot-ccsubspace*[*simp*]: $\langle (\sum x \in X. \perp) = (\perp :: - \text{ ccsubspace}) \rangle$
 $\langle \text{proof} \rangle$

7.7 Conjugate space

```

typedef 'a conjugate-space = UNIV :: 'a set
morphisms from-conjugate-space to-conjugate-space <proof>
setup-lifting type-definition-conjugate-space

instantiation conjugate-space :: (complex-vector) complex-vector begin
lift-definition scaleC-conjugate-space :: <complex  $\Rightarrow$  'a conjugate-space  $\Rightarrow$  'a conjugate-space> is < $\lambda c x. \text{cnj } c *_C x$ ><proof>
lift-definition scaleR-conjugate-space :: <real  $\Rightarrow$  'a conjugate-space  $\Rightarrow$  'a conjugate-space> is < $\lambda r x. r *_R x$ ><proof>
lift-definition plus-conjugate-space :: 'a conjugate-space  $\Rightarrow$  'a conjugate-space  $\Rightarrow$  'a conjugate-space is (+)<proof>
lift-definition uminus-conjugate-space :: 'a conjugate-space  $\Rightarrow$  'a conjugate-space is < $\lambda x. -x$ ><proof>
lift-definition zero-conjugate-space :: 'a conjugate-space is 0<proof>
lift-definition minus-conjugate-space :: 'a conjugate-space  $\Rightarrow$  'a conjugate-space  $\Rightarrow$  'a conjugate-space is (-)<proof>
instance
  <proof>
end

instantiation conjugate-space :: (complex-normed-vector) complex-normed-vector
begin
lift-definition sgn-conjugate-space :: 'a conjugate-space  $\Rightarrow$  'a conjugate-space is sgn<proof>
lift-definition norm-conjugate-space :: 'a conjugate-space  $\Rightarrow$  real is norm<proof>
lift-definition dist-conjugate-space :: 'a conjugate-space  $\Rightarrow$  'a conjugate-space  $\Rightarrow$  real is dist<proof>
lift-definition uniformity-conjugate-space :: ('a conjugate-space  $\times$  'a conjugate-space) filter is uniformity<proof>
lift-definition open-conjugate-space :: 'a conjugate-space set  $\Rightarrow$  bool is open<proof>
instance
  <proof>
end

instantiation conjugate-space :: (cbanach) cbanach begin
instance
  <proof>
end

lemma bounded-antilinear-to-conjugate-space[simp]: <bounded-antilinear to-conjugate-space>
  <proof>

lemma bounded-antilinear-from-conjugate-space[simp]: <bounded-antilinear from-conjugate-space>
  <proof>

lemma antilinear-to-conjugate-space[simp]: <antilinear to-conjugate-space>
  <proof>

```

lemma *antilinear-from-conjugate-space*[simp]: $\langle \text{antilinear from-conjugate-space} \rangle$
 $\langle \text{proof} \rangle$

lemma *cspan-to-conjugate-space*[simp]: $\text{cspan } (\text{to-conjugate-space } 'X) = \text{to-conjugate-space } ' \text{cspan } X$
 $\langle \text{proof} \rangle$

lemma *surj-to-conjugate-space*[simp]: $\text{surj to-conjugate-space}$
 $\langle \text{proof} \rangle$

lemmas *has-derivative-scaleC*[simp, derivative-intros] =
 $\text{bounded-bilinear.FDERIV}[OF \text{ bounded-cbilinear-scaleC}[THEN \text{ bounded-cbilinear.bounded-bilinear}]]$

lemma *norm-to-conjugate-space*[simp]: $\langle \text{norm } (\text{to-conjugate-space } x) = \text{norm } x \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-from-conjugate-space*[simp]: $\langle \text{norm } (\text{from-conjugate-space } x) = \text{norm } x \rangle$
 $\langle \text{proof} \rangle$

lemma *closure-to-conjugate-space*: $\langle \text{closure } (\text{to-conjugate-space } 'X) = \text{to-conjugate-space } ' \text{closure } X \rangle$
 $\langle \text{proof} \rangle$

lemma *closure-from-conjugate-space*: $\langle \text{closure } (\text{from-conjugate-space } 'X) = \text{from-conjugate-space } ' \text{closure } X \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-antilinear-eq-on*:
fixes $A B :: 'a::\text{complex-normed-vector} \Rightarrow 'b::\text{complex-normed-vector}$
assumes $\langle \text{bounded-antilinear } A \rangle$ **and** $\langle \text{bounded-antilinear } B \rangle$ **and**
 $\text{eq: } \langle \bigwedge x. x \in G \implies A x = B x \rangle$ **and** $t: \langle t \in \text{closure } (\text{cspan } G) \rangle$
shows $\langle A t = B t \rangle$
 $\langle \text{proof} \rangle$

lemma *to-conjugate-space-0*[simp]: $\langle \text{to-conjugate-space } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *from-conjugate-space-0*[simp]: $\langle \text{from-conjugate-space } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *antilinear-eq-0-on-span*:
assumes $\langle \text{antilinear } f \rangle$
and $\langle \bigwedge x. x \in b \implies f x = 0 \rangle$
and $\langle x \in \text{cspan } b \rangle$
shows $\langle f x = 0 \rangle$
 $\langle \text{proof} \rangle$

7.8 Separating sets

lemma *separating-set-bounded-clinear-dense*:

assumes $\langle cspan\ S = \top \rangle$
shows $\langle separating\text{-}set\ bounded\text{-}clinear\ S \rangle$
 $\langle proof \rangle$

lemma *separating-set-bounded-cbilinear-nested*:

assumes $\langle separating\text{-}set\ (bounded\text{-}clinear :: (- \Rightarrow 'e :: complex\text{-}normed\text{-}vector) \Rightarrow -) ((\lambda(x, y). h\ x\ y) \text{ ` } (UNIV \times UNIV)) \rangle$
assumes $\langle bounded\text{-}cbilinear\ h \rangle$
assumes $\langle separating\text{-}set\ (bounded\text{-}clinear :: (- \Rightarrow 'e) \Rightarrow -) A \rangle$
assumes $\langle separating\text{-}set\ (bounded\text{-}clinear :: (- \Rightarrow 'e) \Rightarrow -) B \rangle$
shows $\langle separating\text{-}set\ (bounded\text{-}clinear :: (- \Rightarrow 'e) \Rightarrow -) ((\lambda(x, y). h\ x\ y) \text{ ` } (A \times B)) \rangle$
 $\langle proof \rangle$

lemma *separating-set-bounded-clinear-antilinear*:

assumes $\langle separating\text{-}set\ (bounded\text{-}clinear :: (- \Rightarrow 'e :: complex\text{-}normed\text{-}vector\ conjugate\text{-}space) \Rightarrow -) A \rangle$
shows $\langle separating\text{-}set\ (bounded\text{-}antilinear :: (- \Rightarrow 'e) \Rightarrow -) A \rangle$
 $\langle proof \rangle$

lemma *separating-set-bounded-sesquilinear-nested*:

assumes $\langle separating\text{-}set\ (bounded\text{-}clinear :: (- \Rightarrow 'e :: complex\text{-}normed\text{-}vector) \Rightarrow -) ((\lambda(x, y). h\ x\ y) \text{ ` } (UNIV \times UNIV)) \rangle$
assumes $\langle bounded\text{-}sesquilinear\ h \rangle$
assumes *sep-A*: $\langle separating\text{-}set\ (bounded\text{-}clinear :: (- \Rightarrow 'e\ conjugate\text{-}space) \Rightarrow -) A \rangle$
assumes *sep-B*: $\langle separating\text{-}set\ (bounded\text{-}clinear :: (- \Rightarrow 'e) \Rightarrow -) B \rangle$
shows $\langle separating\text{-}set\ (bounded\text{-}clinear :: (- \Rightarrow 'e) \Rightarrow -) ((\lambda(x, y). h\ x\ y) \text{ ` } (A \times B)) \rangle$
 $\langle proof \rangle$

lemma *separating-set-clinear-cspan*:

assumes $\langle cspan\ S = UNIV \rangle$
shows $\langle separating\text{-}set\ clinear\ S \rangle$
 $\langle proof \rangle$

7.9 Product is a Complex Vector Space

instantiation *prod* :: $(complex\text{-}vector, complex\text{-}vector)\ complex\text{-}vector$
begin

definition *scaleC-prod-def*:

$scaleC\ r\ A = (scaleC\ r\ (fst\ A), scaleC\ r\ (snd\ A))$

lemma *fst-scaleC [simp]*: $fst\ (scaleC\ r\ A) = scaleC\ r\ (fst\ A)$

$\langle \text{proof} \rangle$

lemma *snd-scaleC [simp]*: $\text{snd} (\text{scaleC } r \ A) = \text{scaleC } r (\text{snd } A)$
 $\langle \text{proof} \rangle$

proposition *scaleC-Pair [simp]*: $\text{scaleC } r \ (a, b) = (\text{scaleC } r \ a, \text{scaleC } r \ b)$
 $\langle \text{proof} \rangle$

instance
 $\langle \text{proof} \rangle$

end

lemma *module-prod-scale-eq-scaleC*: $\text{module-prod.scale } (*_C) (*_C) = \text{scaleC}$
 $\langle \text{proof} \rangle$

interpretation *complex-vector?: vector-space-prod scaleC:: $\Rightarrow$ $\Rightarrow$ 'a::complex-vector scaleC:: $\Rightarrow$ $\Rightarrow$ 'b::complex-vector*
rewrites $\text{scale} = ((*_C)::\Rightarrow\Rightarrow('a \times 'b))$
and $\text{module.dependent } (*_C) = \text{cdependent}$
and $\text{module.representation } (*_C) = \text{crepresentation}$
and $\text{module.subspace } (*_C) = \text{csubspace}$
and $\text{module.span } (*_C) = \text{cspan}$
and $\text{vector-space.extend-basis } (*_C) = \text{cextend-basis}$
and $\text{vector-space.dim } (*_C) = \text{cdim}$
and $\text{Vector-Spaces.linear } (*_C) (*_C) = \text{clinear}$
 $\langle \text{proof} \rangle$

instance *prod :: (complex-normed-vector, complex-normed-vector) complex-normed-vector*
 $\langle \text{proof} \rangle$

lemma *cspan-Times*: $\langle \text{cspan } (S \times T) = \text{cspan } S \times \text{cspan } T \rangle$ **if** $\langle 0 \in S \rangle$ **and** $\langle 0 \in T \rangle$
 $\langle \text{proof} \rangle$

lemma *onorm-case-prod-plus*: $\langle \text{onorm } (\text{case-prod plus} :: - \Rightarrow 'a::\{\text{real-normed-vector, not-singleton}\}) = \text{sqrt } 2 \rangle$
 $\langle \text{proof} \rangle$

7.10 Copying existing theorems into sublocales

context *bounded-clinear* **begin**
interpretation *bounded-linear f* $\langle \text{proof} \rangle$
lemmas *continuous* = real.continuous
lemmas *uniform-limit* = $\text{real.uniform-limit}$
lemmas *Cauchy* = real.Cauchy
end

```

context bounded-antilinear begin
interpretation bounded-linear f  $\langle \text{proof} \rangle$ 
lemmas continuous = real.continuous
lemmas uniform-limit = real.uniform-limit
end

```

```

context bounded-cbilinear begin
interpretation bounded-bilinear prod  $\langle \text{proof} \rangle$ 
lemmas tendsto = real.tendsto
lemmas isCont = real.isCont
lemmas scaleR-right = real.scaleR-right
lemmas scaleR-left = real.scaleR-left
end

```

```

context bounded-sesquilinear begin
interpretation bounded-bilinear prod  $\langle \text{proof} \rangle$ 
lemmas tendsto = real.tendsto
lemmas isCont = real.isCont
lemmas scaleR-right = real.scaleR-right
lemmas scaleR-left = real.scaleR-left
end

```

```

lemmas tendsto-scaleC [tendsto-intros] =
  bounded-cbilinear.tendsto [OF bounded-cbilinear-scaleC]

```

```

unbundle no lattice-syntax

```

```

end

```

8 *Complex-Inner-Product0* – Inner Product Spaces and Gradient Derivative

```

theory Complex-Inner-Product0
imports
  Complex-Main Complex-Vector-Spaces
  HOL-Analysis.Inner-Product
  Complex-Bounded-Operators.Extra-Ordered-Fields
begin

```

8.1 Complex inner product spaces

Temporarily relax type constraints for *open*, *uniformity*, *dist*, and *norm*.

$\langle ML \rangle$

```

class complex-inner = complex-vector + sgn-div-norm + dist-norm + uniformity-dist + open-uniformity +

```

```

fixes cinner :: 'a  $\Rightarrow$  'a  $\Rightarrow$  complex
assumes cinner-commute: cinner x y = cnj (cinner y x)
  and cinner-add-left: cinner (x + y) z = cinner x z + cinner y z
  and cinner-scaleC-left [simp]: cinner (scaleC r x) y = (cnj r) * (cinner x y)
  and cinner-ge-zero [simp]: 0  $\leq$  cinner x x
  and cinner-eq-zero-iff [simp]: cinner x x = 0  $\longleftrightarrow$  x = 0
  and norm-eq-sqrt-cinner: norm x = sqrt (cmod (cinner x x))
begin

lemma cinner-zero-left [simp]: cinner 0 x = 0
   $\langle$ proof $\rangle$ 

lemma cinner-minus-left [simp]: cinner (- x) y = - cinner x y
   $\langle$ proof $\rangle$ 

lemma cinner-diff-left: cinner (x - y) z = cinner x z - cinner y z
   $\langle$ proof $\rangle$ 

lemma cinner-sum-left: cinner ( $\sum_{x \in A} f$  x) y = ( $\sum_{x \in A} cinner$  (f x) y)
   $\langle$ proof $\rangle$ 

lemma call-zero-iff [simp]: ( $\forall u. cinner$  x u = 0)  $\longleftrightarrow$  (x = 0)
   $\langle$ proof $\rangle$ 

Transfer distributivity rules to right argument.

lemma cinner-add-right: cinner x (y + z) = cinner x y + cinner x z
   $\langle$ proof $\rangle$ 

lemma cinner-scaleC-right [simp]: cinner x (scaleC r y) = r * (cinner x y)
   $\langle$ proof $\rangle$ 

lemma cinner-zero-right [simp]: cinner x 0 = 0
   $\langle$ proof $\rangle$ 

lemma cinner-minus-right [simp]: cinner x (- y) = - cinner x y
   $\langle$ proof $\rangle$ 

lemma cinner-diff-right: cinner x (y - z) = cinner x y - cinner x z
   $\langle$ proof $\rangle$ 

lemma cinner-sum-right: cinner x ( $\sum_{y \in A} f$  y) = ( $\sum_{y \in A} cinner$  x (f y))
   $\langle$ proof $\rangle$ 

lemmas cinner-add [algebra-simps] = cinner-add-left cinner-add-right
lemmas cinner-diff [algebra-simps] = cinner-diff-left cinner-diff-right
lemmas cinner-scaleC = cinner-scaleC-left cinner-scaleC-right

```


lemma *cinner-gt-zero-iff* [simp]: $0 < \text{cinner } x \, x \longleftrightarrow x \neq 0$
 <proof>

lemma *power2-norm-eq-cinner*:
shows $(\text{complex-of-real } (\text{norm } x))^2 = (\text{cinner } x \, x)$
 <proof>

lemma *power2-norm-eq-cinner'*:
shows $(\text{norm } x)^2 = \text{Re } (\text{cinner } x \, x)$
 <proof>

Identities involving real multiplication and division.

lemma *cinner-mult-left*: $\text{cinner } (\text{of-complex } m * a) \, b = \text{cnj } m * (\text{cinner } a \, b)$
 <proof>

lemma *cinner-mult-right*: $\text{cinner } a \, (\text{of-complex } m * b) = m * (\text{cinner } a \, b)$
 <proof>

lemma *cinner-mult-left'*: $\text{cinner } (a * \text{of-complex } m) \, b = \text{cnj } m * (\text{cinner } a \, b)$
 <proof>

lemma *cinner-mult-right'*: $\text{cinner } a \, (b * \text{of-complex } m) = (\text{cinner } a \, b) * m$
 <proof>

lemma *Cauchy-Schwarz-ineq*:
 $(\text{cinner } x \, y) * (\text{cinner } y \, x) \leq \text{cinner } x \, x * \text{cinner } y \, y$
 <proof>

lemma *Cauchy-Schwarz-ineq2*:
shows $\text{norm } (\text{cinner } x \, y) \leq \text{norm } x * \text{norm } y$
 <proof>

subclass *complex-normed-vector*
 <proof>

end

lemma *csquare-continuous*:

fixes $e :: \text{real}$
shows $e > 0 \implies \exists d. 0 < d \wedge (\forall y. \text{cmod } (y - x) < d \longrightarrow \text{cmod } (y * y - x * x) < e)$
 $\langle \text{proof} \rangle$

lemma *cnorm-le*: $\text{norm } x \leq \text{norm } y \longleftrightarrow \text{cinner } x \ x \leq \text{cinner } y \ y$
 $\langle \text{proof} \rangle$

lemma *cnorm-lt*: $\text{norm } x < \text{norm } y \longleftrightarrow \text{cinner } x \ x < \text{cinner } y \ y$
 $\langle \text{proof} \rangle$

lemma *cnorm-eq*: $\text{norm } x = \text{norm } y \longleftrightarrow \text{cinner } x \ x = \text{cinner } y \ y$
 $\langle \text{proof} \rangle$

lemma *cnorm-eq-1*: $\text{norm } x = 1 \longleftrightarrow \text{cinner } x \ x = 1$
 $\langle \text{proof} \rangle$

lemma *cinner-divide-left*:
fixes $a :: 'a :: \{\text{complex-inner}, \text{complex-div-algebra}\}$
shows $\text{cinner } (a / \text{of-complex } m) \ b = (\text{cinner } a \ b) / \text{cnj } m$
 $\langle \text{proof} \rangle$

lemma *cinner-divide-right*:
fixes $a :: 'a :: \{\text{complex-inner}, \text{complex-div-algebra}\}$
shows $\text{cinner } a \ (b / \text{of-complex } m) = (\text{cinner } a \ b) / m$
 $\langle \text{proof} \rangle$

Re-enable constraints for *open*, *uniformity*, *dist*, and *norm*.
 $\langle ML \rangle$

lemma *bounded-sesquilinear-cinner*:
 $\text{bounded-sesquilinear } (\text{cinner} :: 'a :: \text{complex-inner} \Rightarrow 'a \Rightarrow \text{complex})$
 $\langle \text{proof} \rangle$

lemmas *tendsto-cinner* [*tendsto-intros*] =
 $\text{bounded-bilinear.tendsto } [OF \ \text{bounded-sesquilinear-cinner} [THEN \ \text{bounded-sesquilinear.bounded-bilinear}]]$

lemmas *isCont-cinner* [*simp*] =
 $\text{bounded-bilinear.isCont } [OF \ \text{bounded-sesquilinear-cinner} [THEN \ \text{bounded-sesquilinear.bounded-bilinear}]]$

lemmas *has-derivative-cinner* [*derivative-intros*] =
 $\text{bounded-bilinear.FDERIV } [OF \ \text{bounded-sesquilinear-cinner} [THEN \ \text{bounded-sesquilinear.bounded-bilinear}]]$

lemmas *bounded-antilinear-cinner-left* =
 $\text{bounded-sesquilinear.bounded-antilinear-left } [OF \ \text{bounded-sesquilinear-cinner}]$

lemmas *bounded-clinear-cinner-right* =
 $\text{bounded-sesquilinear.bounded-clinear-right } [OF \ \text{bounded-sesquilinear-cinner}]$

lemmas *bounded-antilinear-cinner-left-comp* = *bounded-antilinear-cinner-left*[*THEN*
bounded-antilinear-o-bounded-clinear]

lemmas *bounded-clinear-cinner-right-comp* = *bounded-clinear-cinner-right*[*THEN*
bounded-clinear-compose]

lemmas *has-derivative-cinner-right* [*derivative-intros*] =
bounded-linear.has-derivative [*OF* *bounded-clinear-cinner-right*[*THEN* *bounded-clinear.bounded-linear*]]

lemmas *has-derivative-cinner-left* [*derivative-intros*] =
bounded-linear.has-derivative [*OF* *bounded-antilinear-cinner-left*[*THEN* *bounded-antilinear.bounded-linear*]]

lemma *differentiable-cinner* [*simp*]:
 $f \text{ differentiable (at } x \text{ within } s) \implies g \text{ differentiable at } x \text{ within } s \implies (\lambda x. \text{ cinner } (f \ x) \ (g \ x)) \text{ differentiable at } x \text{ within } s$
<proof>

8.2 Class instances

instantiation *complex* :: *complex-inner*
begin

definition *cinner-complex-def* [*simp*]: *cinner* $x \ y = \text{cnj } x * y$

instance
<proof>

end

lemma
shows *complex-inner-1-left*[*simp*]: *cinner* 1 $x = x$
and *complex-inner-1-right*[*simp*]: *cinner* x 1 = *cnj* x
<proof>

lemma *cdot-square-norm*: *cinner* $x \ x = \text{complex-of-real } ((\text{norm } x)^2)$
<proof>

lemma *cnorm-eq-square*: $\text{norm } x = a \iff 0 \leq a \wedge \text{cinner } x \ x = \text{complex-of-real } (a^2)$
<proof>

lemma *cnorm-le-square*: $\text{norm } x \leq a \iff 0 \leq a \wedge \text{cinner } x \ x \leq \text{complex-of-real } (a^2)$
<proof>

lemma *cnorm-ge-square*: $\text{norm } x \geq a \iff a \leq 0 \vee \text{cinner } x \ x \geq \text{complex-of-real } (a^2)$

(a^2)
 $\langle \text{proof} \rangle$

lemma *norm-lt-square*: $\text{norm } x < a \iff 0 < a \wedge \text{cinner } x \ x < \text{complex-of-real } (a^2)$
 $\langle \text{proof} \rangle$

lemma *norm-gt-square*: $\text{norm } x > a \iff a < 0 \vee \text{cinner } x \ x > \text{complex-of-real } (a^2)$
 $\langle \text{proof} \rangle$

Dot product in terms of the norm rather than conversely.

lemmas *cinner-simps* = *cinner-add-left cinner-add-right cinner-diff-right cinner-diff-left cinner-scaleC-left cinner-scaleC-right*

lemma *cdot-norm*: $\text{cinner } x \ y = ((\text{norm } (x+y))^2 - (\text{norm } (x-y))^2 - i * (\text{norm } (x + i *_C y))^2 + i * (\text{norm } (x - i *_C y))^2) / 4$
 $\langle \text{proof} \rangle$

lemma *of-complex-inner-1* [simp]:
 $\text{cinner } (\text{of-complex } x) \ (1 :: 'a :: \{\text{complex-inner}, \text{complex-normed-algebra-1}\}) = \text{cnj } x$
 $\langle \text{proof} \rangle$

lemma *summable-of-complex-iff*:
 $\text{summable } (\lambda x. \text{of-complex } (f x) :: 'a :: \{\text{complex-normed-algebra-1}, \text{complex-inner}\}) \iff \text{summable } f$
 $\langle \text{proof} \rangle$

8.3 Gradient derivative

definition
 $\text{cgderiv} :: ['a :: \text{complex-inner} \Rightarrow \text{complex}, 'a, 'a] \Rightarrow \text{bool}$
 $(\langle \text{cGDERIV } (-) / (-) / :> (-) \rangle [1000, 1000, 60] 60)$
where

$\text{cGDERIV } f \ x :> D \iff \text{FDERIV } f \ x :> \text{cinner } D$

lemma *cgderiv-deriv* [simp]: $\text{cGDERIV } f \ x :> D \iff \text{DERIV } f \ x :> \text{cnj } D$
 $\langle \text{proof} \rangle$

lemma *cGDERIV-DERIV-compose*:
assumes $\text{cGDERIV } f \ x :> df$ **and** $\text{DERIV } g \ (f \ x) :> \text{cnj } dg$
shows $\text{cGDERIV } (\lambda x. g \ (f \ x)) \ x :> \text{scaleC } dg \ df$
 $\langle \text{proof} \rangle$

lemma *cGDERIV-subst*: $\llbracket cGDERIV\ f\ x\ :\>\ df;\ df = d \rrbracket \implies cGDERIV\ f\ x\ :\>\ d$
 $\langle proof \rangle$

lemma *cGDERIV-const*: $cGDERIV\ (\lambda x. k)\ x\ :\>\ 0$
 $\langle proof \rangle$

lemma *cGDERIV-add*:
 $\llbracket cGDERIV\ f\ x\ :\>\ df;\ cGDERIV\ g\ x\ :\>\ dg \rrbracket$
 $\implies cGDERIV\ (\lambda x. f\ x + g\ x)\ x\ :\>\ df + dg$
 $\langle proof \rangle$

lemma *cGDERIV-minus*:
 $cGDERIV\ f\ x\ :\>\ df \implies cGDERIV\ (\lambda x. - f\ x)\ x\ :\>\ - df$
 $\langle proof \rangle$

lemma *cGDERIV-diff*:
 $\llbracket cGDERIV\ f\ x\ :\>\ df;\ cGDERIV\ g\ x\ :\>\ dg \rrbracket$
 $\implies cGDERIV\ (\lambda x. f\ x - g\ x)\ x\ :\>\ df - dg$
 $\langle proof \rangle$

lemma *cGDERIV-scaleC*:
 $\llbracket DERIV\ f\ x\ :\>\ df;\ cGDERIV\ g\ x\ :\>\ dg \rrbracket$
 $\implies cGDERIV\ (\lambda x. scaleC\ (f\ x)\ (g\ x))\ x$
 $\quad :\>\ (scaleC\ (cnj\ (f\ x))\ dg + scaleC\ (cnj\ df)\ (cnj\ (g\ x)))$
 $\langle proof \rangle$

lemma *GDERIV-mult*:
 $\llbracket cGDERIV\ f\ x\ :\>\ df;\ cGDERIV\ g\ x\ :\>\ dg \rrbracket$
 $\implies cGDERIV\ (\lambda x. f\ x * g\ x)\ x\ :\>\ cnj\ (f\ x) *_C dg + cnj\ (g\ x) *_C df$
 $\langle proof \rangle$

lemma *cGDERIV-inverse*:
 $\llbracket cGDERIV\ f\ x\ :\>\ df;\ f\ x \neq 0 \rrbracket$
 $\implies cGDERIV\ (\lambda x. inverse\ (f\ x))\ x\ :\>\ - cnj\ ((inverse\ (f\ x))^2) *_C df$
 $\langle proof \rangle$

lemma *has-derivative-norm*[*derivative-intros*]:
fixes $x :: 'a::complex-inner$
assumes $x \neq 0$
shows $(norm\ has-derivative\ (\lambda h. Re\ (cinner\ (sgn\ x)\ h)))\ (at\ x)$
thm *has-derivative-norm*
 $\langle proof \rangle$

```

bundle cinner-syntax
begin
notation cinner (infix  $\langle \cdot_C \rangle$  70)
end

end

```

9 Complex-Inner-Product – Complex Inner Product Spaces

```

theory Complex-Inner-Product
imports
  Complex-Inner-Product0
begin

```

9.1 Complex inner product spaces

```

unbundle cinner-syntax

```

```

lemma cinner-real: cinner  $x$   $x \in \mathbb{R}$ 
   $\langle proof \rangle$ 

```

```

lemmas cinner-commute' [simp] = cinner-commute[symmetric]

```

```

lemma (in complex-inner) cinner-eq-flip:  $\langle (cinner\ x\ y = cinner\ z\ w) \longleftrightarrow (cinner\ y\ x = cinner\ w\ z) \rangle$ 
   $\langle proof \rangle$ 

```

```

lemma Im-cinner-x-x[simp]:  $Im\ (x \cdot_C x) = 0$ 
   $\langle proof \rangle$ 

```

```

lemma of-complex-inner-1' [simp]:
  cinner (1 :: 'a :: {complex-inner, complex-normed-algebra-1}) (of-complex  $x$ ) =
   $x$ 
   $\langle proof \rangle$ 

```

```

class hilbert-space = complex-inner + complete-space
begin
subclass cbanach  $\langle proof \rangle$ 
end

```

```

instantiation complex :: hilbert-space begin
instance  $\langle proof \rangle$ 
end

```

9.2 Misc facts

lemma *cinner-scaleR-left* [simp]: $\text{cinner} (\text{scaleR } r \ x) \ y = \text{of-real } r * (\text{cinner } x \ y)$
 ⟨proof⟩

lemma *cinner-scaleR-right* [simp]: $\text{cinner } x (\text{scaleR } r \ y) = \text{of-real } r * (\text{cinner } x \ y)$
 ⟨proof⟩

This is a useful rule for establishing the equality of vectors

lemma *cinner-extensionality*:
assumes $\langle \bigwedge \gamma. \gamma \cdot_C \psi = \gamma \cdot_C \varphi \rangle$
shows $\langle \psi = \varphi \rangle$
 ⟨proof⟩

lemma *polar-identity*:
includes *norm-syntax*
shows $\langle \|x + y\|^2 = \|x\|^2 + \|y\|^2 + 2 * \text{Re } (x \cdot_C y) \rangle$
 — Shown in the proof of Corollary 1.5 in [1]
 ⟨proof⟩

lemma *polar-identity-minus*:
includes *norm-syntax*
shows $\langle \|x - y\|^2 = \|x\|^2 + \|y\|^2 - 2 * \text{Re } (x \cdot_C y) \rangle$
 ⟨proof⟩

proposition *parallelogram-law*:
includes *norm-syntax*
fixes $x \ y :: 'a::\text{complex-inner}$
shows $\langle \|x+y\|^2 + \|x-y\|^2 = 2*(\|x\|^2 + \|y\|^2) \rangle$
 — Shown in the proof of Theorem 2.3 in [1]
 ⟨proof⟩

theorem *pythagorean-theorem*:
includes *norm-syntax*
shows $\langle (x \cdot_C y) = 0 \implies \|x + y\|^2 = \|x\|^2 + \|y\|^2 \rangle$
 — Shown in the proof of Theorem 2.2 in [1]
 ⟨proof⟩

lemma *pythagorean-theorem-sum*:
assumes $q1: \bigwedge a \ a'. a \in t \implies a' \in t \implies a \neq a' \implies f \ a \cdot_C f \ a' = 0$
and $q2: \text{finite } t$
shows $(\text{norm } (\sum a \in t. f \ a))^2 = (\sum a \in t. (\text{norm } (f \ a))^2)$
 ⟨proof⟩

lemma *Cauchy-cinner-Cauchy*:
fixes $x \ y :: \langle \text{nat} \Rightarrow 'a::\text{complex-inner} \rangle$
assumes $a1: \langle \text{Cauchy } x \rangle$ **and** $a2: \langle \text{Cauchy } y \rangle$
shows $\langle \text{Cauchy } (\lambda n. x \ n \cdot_C y \ n) \rangle$

$\langle \text{proof} \rangle$

lemma *cinner-sup-norm*: $\langle \text{norm } \psi = (\text{SUP } \varphi. \text{cmod } (\text{cinner } \varphi \ \psi) / \text{norm } \varphi) \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-sup-onorm*:
fixes $A :: \langle 'a::\{\text{real-normed-vector}, \text{not-singleton}\} \Rightarrow 'b::\text{complex-inner} \rangle$
assumes $\langle \text{bounded-linear } A \rangle$
shows $\langle \text{onorm } A = (\text{SUP } (\psi, \varphi). \text{cmod } (\text{cinner } \psi \ (A \ \varphi)) / (\text{norm } \psi * \text{norm } \varphi)) \rangle$
 $\langle \text{proof} \rangle$

lemma *sum-cinner*:
fixes $f :: 'a \Rightarrow 'b::\text{complex-inner}$
shows $\langle \text{cinner } (\text{sum } f \ A) \ (\text{sum } g \ B) = (\sum i \in A. \sum j \in B. \text{cinner } (f \ i) \ (g \ j)) \rangle$
 $\langle \text{proof} \rangle$

lemma *Cauchy-cinner-product-summable'*:
fixes $a \ b :: \text{nat} \Rightarrow 'a::\text{complex-inner}$
shows $\langle (\lambda(x, y). \text{cinner } (a \ x) \ (b \ y)) \text{ summable-on } \text{UNIV} \longleftrightarrow (\lambda(x, y). \text{cinner } (a \ y) \ (b \ (x - y))) \text{ summable-on } \{(k, i). i \leq k\} \rangle$
 $\langle \text{proof} \rangle$

instantiation *prod* :: $(\text{complex-inner}, \text{complex-inner}) \text{ complex-inner}$
begin

definition *cinner-prod-def*:
 $\text{cinner } x \ y = \text{cinner } (\text{fst } x) \ (\text{fst } y) + \text{cinner } (\text{snd } x) \ (\text{snd } y)$

instance
 $\langle \text{proof} \rangle$

end

lemma *sgn-cinner[simp]*: $\langle \text{sgn } \psi \cdot_C \psi = \text{norm } \psi \rangle$
 $\langle \text{proof} \rangle$

instance *prod* :: $(\text{chilbert-space}, \text{chilbert-space}) \text{ chilbert-space} \langle \text{proof} \rangle$

9.3 Orthogonality

definition *orthogonal-complement* $S = \{x \mid x. \forall y \in S. \text{cinner } x \ y = 0\}$

lemma *orthogonal-complement-orthoI*:
 $\langle x \in \text{orthogonal-complement } M \implies y \in M \implies x \cdot_C y = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complement-orthoI'*:

$\langle x \in M \implies y \in \text{orthogonal-complement } M \implies x \cdot_C y = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complementI*:
 $\langle (\bigwedge x. x \in M \implies y \cdot_C x = 0) \implies y \in \text{orthogonal-complement } M \rangle$
 $\langle \text{proof} \rangle$

abbreviation *is-orthogonal*:: $\langle 'a::\text{complex-inner} \Rightarrow 'a \Rightarrow \text{bool} \rangle$ **where**
 $\langle \text{is-orthogonal } x \ y \equiv x \cdot_C y = 0 \rangle$

bundle *orthogonal-syntax*
begin
notation *is-orthogonal* (**infixl** $\langle \perp \rangle$ 69)
end

lemma *is-orthogonal-sym*: *is-orthogonal* $\psi \ \varphi = \text{is-orthogonal } \varphi \ \psi$
 $\langle \text{proof} \rangle$

lemma *is-orthogonal-sgn-right[simp]*: $\langle \text{is-orthogonal } e \ (\text{sgn } f) \longleftrightarrow \text{is-orthogonal } e \ f \rangle$
 $\langle \text{proof} \rangle$

lemma *is-orthogonal-sgn-left[simp]*: $\langle \text{is-orthogonal } (\text{sgn } e) \ f \longleftrightarrow \text{is-orthogonal } e \ f \rangle$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complement-closed-subspace[simp]*:
closed-csubspace (*orthogonal-complement* A)
for $A :: \langle 'a::\text{complex-inner} \rangle \text{ set}$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complement-zero-intersection*:
assumes $0 \in M$
shows $\langle M \cap (\text{orthogonal-complement } M) = \{0\} \rangle$
 $\langle \text{proof} \rangle$

lemma *is-orthogonal-closure-cspan*:
assumes $\bigwedge x \ y. x \in X \implies y \in Y \implies \text{is-orthogonal } x \ y$
assumes $\langle x \in \text{closure } (\text{cspan } X) \rangle \langle y \in \text{closure } (\text{cspan } Y) \rangle$
shows *is-orthogonal* $x \ y$
 $\langle \text{proof} \rangle$

instantiation *ccsubspace* :: $(\text{complex-inner}) \ \text{uminus}$
begin
lift-definition *uminus-ccsubspace*:: $\langle 'a \ \text{ccsubspace} \Rightarrow 'a \ \text{ccsubspace} \rangle$
is $\langle \text{orthogonal-complement} \rangle$
 $\langle \text{proof} \rangle$

instance $\langle proof \rangle$
end

lemma *orthocomplement-top*[simp]: $\langle - \text{ top} = (\text{bot} :: 'a::\text{complex-inner ccspace}) \rangle$
 — For $'a$ of sort *chilbert-space*, this is covered by *orthocomplemented-lattice-class.compl-top-eq* already. But here we give it a wider sort.
 $\langle proof \rangle$

instantiation *ccspace* :: (*complex-inner*) *minus* **begin**
lift-definition *minus-ccspace* :: $'a \text{ ccspace} \Rightarrow 'a \text{ ccspace} \Rightarrow 'a \text{ ccspace}$
is $\lambda A B. A \cap (\text{orthogonal-complement } B)$
 $\langle proof \rangle$
instance $\langle proof \rangle$
end

definition *is-ortho-set* :: $'a::\text{complex-inner set} \Rightarrow \text{bool}$ **where**
 — Orthogonal set
 $\langle is-ortho-set \ S \longleftrightarrow (\forall x \in S. \forall y \in S. x \neq y \longrightarrow (x \cdot_C y) = 0) \wedge 0 \notin S \rangle$

definition *is-onb* :: $'a::\text{complex-inner set} \Rightarrow \text{bool}$ **where**
 — Orthonormal basis
 $\langle is-onb \ E \longleftrightarrow is-ortho-set \ E \wedge (\forall b \in E. \text{norm } b = 1) \wedge \text{ccspan } E = \text{top} \rangle$

lemma *is-ortho-set-empty*[simp]: *is-ortho-set* $\{\}$
 $\langle proof \rangle$

lemma *is-ortho-set-antimono*: $\langle A \subseteq B \Longrightarrow is-ortho-set \ B \Longrightarrow is-ortho-set \ A \rangle$
 $\langle proof \rangle$

lemma *orthogonal-complement-of-closure*:
fixes $A :: ('a::\text{complex-inner}) \text{ set}$
shows $\text{orthogonal-complement } A = \text{orthogonal-complement } (\text{closure } A)$
 $\langle proof \rangle$

lemma *is-orthogonal-closure*:
assumes $\langle \bigwedge s. s \in S \Longrightarrow is-orthogonal \ a \ s \rangle$
assumes $\langle x \in \text{closure } S \rangle$
shows $\langle is-orthogonal \ a \ x \rangle$
 $\langle proof \rangle$

lemma *is-orthogonal-cspan*:
assumes $a1: \bigwedge s. s \in S \Longrightarrow is-orthogonal \ a \ s$ **and** $a3: x \in \text{cspan } S$
shows $is-orthogonal \ a \ x$
 $\langle proof \rangle$

lemma *ccspan-leq-ortho-ccspan*:
assumes $\bigwedge s \ t. s \in S \Longrightarrow t \in T \Longrightarrow is-orthogonal \ s \ t$

shows $\text{ccspan } S \leq - (\text{ccspan } T)$
 $\langle \text{proof} \rangle$

lemma *double-orthogonal-complement-increasing[simp]*:
shows $M \subseteq \text{orthogonal-complement } (\text{orthogonal-complement } M)$
 $\langle \text{proof} \rangle$

lemma *orthonormal-basis-of-cspan*:
fixes $S :: 'a :: \text{complex-inner set}$
assumes *finite S*
shows $\exists A. \text{is-ortho-set } A \wedge (\forall x \in A. \text{norm } x = 1) \wedge \text{cspan } A = \text{cspan } S \wedge \text{finite } A$
 $\langle \text{proof} \rangle$

lemma *is-ortho-set-cindependent*:
assumes *is-ortho-set A*
shows *cindependent A*
 $\langle \text{proof} \rangle$

lemma *onb-expansion-finite*:
includes *norm-syntax*
fixes $T :: 'a :: \{\text{complex-inner}, \text{cfinite-dim}\} \text{ set}$
assumes $a1: \langle \text{cspan } T = \text{UNIV} \rangle$ **and** $a3: \langle \text{is-ortho-set } T \rangle$
and $a4: \langle \bigwedge t. t \in T \implies \|t\| = 1 \rangle$
shows $\langle x = (\sum t \in T. (t \cdot_C x) *_C t) \rangle$
 $\langle \text{proof} \rangle$

lemma *is-ortho-set-singleton[simp]*: $\langle \text{is-ortho-set } \{x\} \longleftrightarrow x \neq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complement-antimono[simp]*:
fixes $A B :: \langle 'a :: \text{complex-inner} \rangle \text{ set}$
assumes $A \supseteq B$
shows $\langle \text{orthogonal-complement } A \subseteq \text{orthogonal-complement } B \rangle$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complement-UNIV[simp]*:
 $\text{orthogonal-complement } \text{UNIV} = \{0\}$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complement-zero[simp]*:
 $\text{orthogonal-complement } \{0\} = \text{UNIV}$
 $\langle \text{proof} \rangle$

lemma *mem-ortho-ccspanI*:
assumes $\langle \bigwedge y. y \in S \implies \text{is-orthogonal } x \ y \rangle$
shows $\langle x \in \text{space-as-set } (- \text{ccspan } S) \rangle$

⟨proof⟩

lemma *cfinite-dim-subspace-has-onb*:

assumes ⟨*cfinite-dim* S ⟩ and ⟨*csubspace* S ⟩

shows ⟨ $\exists B. \text{finite } B \wedge \text{is-ortho-set } B \wedge \text{cspan } B = S \wedge (\forall x \in B. \text{norm } x = 1)$ ⟩

⟨proof⟩

9.4 Projections

lemma *smallest-norm-exists*:

— Theorem 2.5 in [1] (inside the proof)

includes *norm-syntax*

fixes $M :: \langle 'a :: \text{hilbert-space set} \rangle$

assumes $q1: \langle \text{convex } M \rangle$ and $q2: \langle \text{closed } M \rangle$ and $q3: \langle M \neq \{\} \rangle$

shows ⟨ $\exists k. \text{is-arg-min } (\lambda x. \|x\|) (\lambda t. t \in M) k$ ⟩

⟨proof⟩

lemma *smallest-norm-unique*:

— Theorem 2.5 in [1] (inside the proof)

includes *norm-syntax*

fixes $M :: \langle 'a :: \text{complex-inner set} \rangle$

assumes $q1: \langle \text{convex } M \rangle$

assumes $r: \langle \text{is-arg-min } (\lambda x. \|x\|) (\lambda t. t \in M) r \rangle$

assumes $s: \langle \text{is-arg-min } (\lambda x. \|x\|) (\lambda t. t \in M) s \rangle$

shows ⟨ $r = s$ ⟩

⟨proof⟩

theorem *smallest-dist-exists*:

— Theorem 2.5 in [1]

fixes $M :: \langle 'a :: \text{hilbert-space set} \rangle$ and h

assumes $a1: \langle \text{convex } M \rangle$ and $a2: \langle \text{closed } M \rangle$ and $a3: \langle M \neq \{\} \rangle$

shows ⟨ $\exists k. \text{is-arg-min } (\lambda x. \text{dist } x h) (\lambda x. x \in M) k$ ⟩

⟨proof⟩

theorem *smallest-dist-unique*:

— Theorem 2.5 in [1]

fixes $M :: \langle 'a :: \text{complex-inner set} \rangle$ and h

assumes $a1: \langle \text{convex } M \rangle$

assumes ⟨ $\text{is-arg-min } (\lambda x. \text{dist } x h) (\lambda x. x \in M) r$ ⟩

assumes ⟨ $\text{is-arg-min } (\lambda x. \text{dist } x h) (\lambda x. x \in M) s$ ⟩

shows ⟨ $r = s$ ⟩

⟨proof⟩

theorem *smallest-dist-is-ortho*:

fixes $M :: \langle 'a :: \text{complex-inner set} \rangle$ and $h k :: 'a$

assumes $b1: \langle \text{closed-csubspace } M \rangle$

shows ⟨ $(\text{is-arg-min } (\lambda x. \text{dist } x h) (\lambda x. x \in M) k) \longleftrightarrow$
 $h - k \in \text{orthogonal-complement } M \wedge k \in M$ ⟩

⟨proof⟩

include *norm-syntax*
 $\langle \text{proof} \rangle$

corollary *orthog-proj-exists*:
fixes $M :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $\langle \text{closed-csubspace } M \rangle$
shows $\langle \exists k. h - k \in \text{orthogonal-complement } M \wedge k \in M \rangle$
 $\langle \text{proof} \rangle$

corollary *orthog-proj-unique*:
fixes $M :: \langle 'a::\text{complex-inner set} \rangle$
assumes $\langle \text{closed-csubspace } M \rangle$
assumes $\langle h - r \in \text{orthogonal-complement } M \wedge r \in M \rangle$
assumes $\langle h - s \in \text{orthogonal-complement } M \wedge s \in M \rangle$
shows $\langle r = s \rangle$
 $\langle \text{proof} \rangle$

definition *is-projection-on*: $\langle ('a \Rightarrow 'a) \Rightarrow ('a::\text{metric-space}) \text{ set} \Rightarrow \text{bool} \rangle$ **where**
 $\langle \text{is-projection-on } \pi \ M \iff (\forall h. \text{is-arg-min } (\lambda x. \text{dist } x \ h) \ (\lambda x. x \in M) \ (\pi \ h)) \rangle$

lemma *is-projection-on-iff-orthog*:
 $\langle \text{closed-csubspace } M \implies \text{is-projection-on } \pi \ M \iff (\forall h. h - \pi \ h \in \text{orthogonal-complement } M \wedge \pi \ h \in M) \rangle$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-exists*:
fixes $M :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $\langle \text{convex } M \rangle$ **and** $\langle \text{closed } M \rangle$ **and** $\langle M \neq \{\} \rangle$
shows $\exists \pi. \text{is-projection-on } \pi \ M$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-unique*:
fixes $M :: \langle 'a::\text{complex-inner set} \rangle$
assumes $\langle \text{convex } M \rangle$
assumes $\text{is-projection-on } \pi_1 \ M$
assumes $\text{is-projection-on } \pi_2 \ M$
shows $\pi_1 = \pi_2$
 $\langle \text{proof} \rangle$

definition *projection* :: $\langle 'a::\text{metric-space set} \Rightarrow ('a \Rightarrow 'a) \rangle$ **where**
 $\langle \text{projection } M = (\text{SOME } \pi. \text{is-projection-on } \pi \ M) \rangle$

lemma *projection-is-projection-on*:
fixes $M :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $\langle \text{convex } M \rangle$ **and** $\langle \text{closed } M \rangle$ **and** $\langle M \neq \{\} \rangle$
shows $\text{is-projection-on } (\text{projection } M) \ M$
 $\langle \text{proof} \rangle$

lemma *projection-is-projection-on[simp]*:

— Common special case of $\llbracket \text{convex } ?M; \text{closed } ?M; ?M \neq \{\} \rrbracket \implies \text{is-projection-on } (\text{projection } ?M) \text{ } ?M$

fixes $M :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $\langle \text{closed-csubspace } M \rangle$
shows $\text{is-projection-on } (\text{projection } M) \text{ } M$
 $\langle \text{proof} \rangle$

lemma *projection-orthogonal*:
fixes $M :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $\text{closed-csubspace } M$ **and** $\langle m \in M \rangle$
shows $\langle \text{is-orthogonal } (h - \text{projection } M \text{ } h) \text{ } m \rangle$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-in-image*:
assumes $\text{is-projection-on } \pi \text{ } M$
shows $\pi \text{ } h \in M$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-image*:
assumes $\text{is-projection-on } \pi \text{ } M$
shows $\text{range } \pi = M$
 $\langle \text{proof} \rangle$

lemma *projection-in-image[simp]*:
fixes $M :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $\langle \text{convex } M \rangle$ **and** $\langle \text{closed } M \rangle$ **and** $\langle M \neq \{\} \rangle$
shows $\langle \text{projection } M \text{ } h \in M \rangle$
 $\langle \text{proof} \rangle$

lemma *projection-image[simp]*:
fixes $M :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $\langle \text{convex } M \rangle$ **and** $\langle \text{closed } M \rangle$ **and** $\langle M \neq \{\} \rangle$
shows $\langle \text{range } (\text{projection } M) = M \rangle$
 $\langle \text{proof} \rangle$

lemma *projection-eqI'*:
fixes $M :: \langle 'a::\text{complex-inner set} \rangle$
assumes $\langle \text{convex } M \rangle$
assumes $\langle \text{is-projection-on } f \text{ } M \rangle$
shows $\langle \text{projection } M = f \rangle$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-eqI*:
fixes $M :: \langle 'a::\text{complex-inner set} \rangle$
assumes $a1: \langle \text{closed-csubspace } M \rangle$ **and** $a2: \langle h - x \in \text{orthogonal-complement } M \rangle$
and $a3: \langle x \in M \rangle$
and $a4: \langle \text{is-projection-on } \pi \text{ } M \rangle$
shows $\langle \pi \text{ } h = x \rangle$
 $\langle \text{proof} \rangle$

lemma *projection-eqI*:
fixes $M :: \langle 'a::\text{hilbert-space} \rangle \text{ set}$
assumes $\langle \text{closed-csubspace } M \rangle$ **and** $\langle h - x \in \text{orthogonal-complement } M \rangle$ **and**
 $\langle x \in M \rangle$
shows $\langle \text{projection } M \ h = x \rangle$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-fixes-image*:
fixes $M :: \langle 'a::\text{metric-space} \rangle \text{ set}$
assumes $a1: \text{is-projection-on } \pi \ M$ **and** $a3: x \in M$
shows $\pi \ x = x$
 $\langle \text{proof} \rangle$

lemma *projection-fixes-image*:
fixes $M :: \langle 'a::\text{hilbert-space} \rangle \text{ set}$
assumes $\text{closed-csubspace } M$ **and** $x \in M$
shows $\text{projection } M \ x = x$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-closed*:
assumes $\text{cont-}f: \langle \bigwedge x. x \in \text{closure } M \implies \text{isCont } f \ x \rangle$
assumes $\langle \text{is-projection-on } f \ M \rangle$
shows $\langle \text{closed } M \rangle$
 $\langle \text{proof} \rangle$

proposition *is-projection-on-reduces-norm*:
includes *norm-syntax*
fixes $M :: \langle 'a::\text{complex-inner} \rangle \text{ set}$
assumes $\langle \text{is-projection-on } \pi \ M \rangle$ **and** $\langle \text{closed-csubspace } M \rangle$
shows $\langle \| \pi \ h \| \leq \| h \| \rangle$
 $\langle \text{proof} \rangle$

proposition *projection-reduces-norm*:
includes *norm-syntax*
fixes $M :: \langle 'a::\text{hilbert-space} \rangle \text{ set}$
assumes $a1: \text{closed-csubspace } M$
shows $\langle \| \text{projection } M \ h \| \leq \| h \| \rangle$
 $\langle \text{proof} \rangle$

theorem *is-projection-on-bounded-clinear*:
fixes $M :: \langle 'a::\text{complex-inner} \rangle \text{ set}$
assumes $a1: \text{is-projection-on } \pi \ M$ **and** $a2: \text{closed-csubspace } M$
shows $\text{bounded-clinear } \pi$
 $\langle \text{proof} \rangle$

theorem *projection-bounded-clinear*:
fixes $M :: \langle 'a::\text{hilbert-space} \rangle \text{ set}$
assumes $a1: \text{closed-csubspace } M$
shows $\langle \text{bounded-clinear } (\text{projection } M) \rangle$

— Theorem 2.7 in [1]
 $\langle \text{proof} \rangle$

proposition *is-projection-on-idem*:
fixes $M :: \langle 'a::\text{complex-inner} \rangle \text{ set}$
assumes *is-projection-on* π M
shows $\pi (\pi x) = \pi x$
 $\langle \text{proof} \rangle$

proposition *projection-idem*:
fixes $M :: 'a::\text{hilbert-space set}$
assumes $a1: \text{closed-csubspace } M$
shows $\text{projection } M (\text{projection } M x) = \text{projection } M x$
 $\langle \text{proof} \rangle$

proposition *is-projection-on-kernel-is-orthogonal-complement*:
fixes $M :: \langle 'a::\text{complex-inner} \rangle \text{ set}$
assumes $a1: \text{is-projection-on } \pi M$ **and** $a2: \text{closed-csubspace } M$
shows $\pi - \{0\} = \text{orthogonal-complement } M$
 $\langle \text{proof} \rangle$

proposition *projection-kernel-is-orthogonal-complement*:
fixes $M :: \langle 'a::\text{hilbert-space set}$
assumes $\text{closed-csubspace } M$
shows $(\text{projection } M) - \{0\} = (\text{orthogonal-complement } M)$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-id-minus*:
fixes $M :: \langle 'a::\text{complex-inner} \rangle \text{ set}$
assumes *is-proj*: *is-projection-on* π M
and $cc: \text{closed-csubspace } M$
shows *is-projection-on* $(\text{id} - \pi)$ $(\text{orthogonal-complement } M)$
 $\langle \text{proof} \rangle$

Exercise 2 (section 2, chapter I) in [1]

lemma *projection-on-orthogonal-complement[simp]*:
fixes $M :: 'a::\text{hilbert-space set}$
assumes $a1: \text{closed-csubspace } M$
shows $\text{projection } (\text{orthogonal-complement } M) = \text{id} - \text{projection } M$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-zero*:
is-projection-on $(\lambda -. 0)$ $\{0\}$
 $\langle \text{proof} \rangle$

lemma *projection-zero[simp]*:
 $\text{projection } \{0\} = (\lambda -. 0)$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-rank1*:
fixes $t :: \langle 'a::\text{complex-inner} \rangle$
shows $\langle \text{is-projection-on } (\lambda x. ((t \cdot_C x) / (t \cdot_C t)) *_C t) \text{ (cspan } \{t\}) \rangle$
 $\langle \text{proof} \rangle$

lemma *projection-rank1*:
fixes $t x :: \langle 'a::\text{complex-inner} \rangle$
shows $\langle \text{projection (cspan } \{t\}) x = ((t \cdot_C x) / (t \cdot_C t)) *_C t \rangle$
 $\langle \text{proof} \rangle$

9.5 More orthogonal complement

The following lemmas logically fit into the "orthogonality" section but depend on projections for their proofs.

Corollary 2.8 in [1]

theorem *double-orthogonal-complement-id[simp]*:
fixes $M :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $a1: \text{closed-csubspace } M$
shows $\text{orthogonal-complement (orthogonal-complement } M) = M$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complement-antimono-iff[simp]*:
fixes $A B :: \langle 'a::\text{hilbert-space} \rangle \text{ set}$
assumes $\langle \text{closed-csubspace } A \rangle$ **and** $\langle \text{closed-csubspace } B \rangle$
shows $\langle \text{orthogonal-complement } A \subseteq \text{orthogonal-complement } B \iff A \supseteq B \rangle$
 $\langle \text{proof} \rangle$

lemma *de-morgan-orthogonal-complement-plus*:
fixes $A B :: \langle 'a::\text{complex-inner} \rangle \text{ set}$
assumes $\langle 0 \in A \rangle$ **and** $\langle 0 \in B \rangle$
shows $\langle \text{orthogonal-complement } (A +_M B) = \text{orthogonal-complement } A \cap \text{orthogonal-complement } B \rangle$
 $\langle \text{proof} \rangle$

lemma *de-morgan-orthogonal-complement-inter*:
fixes $A B :: 'a::\text{hilbert-space set}$
assumes $a1: \langle \text{closed-csubspace } A \rangle$ **and** $a2: \langle \text{closed-csubspace } B \rangle$
shows $\langle \text{orthogonal-complement } (A \cap B) = \text{orthogonal-complement } A +_M \text{orthogonal-complement } B \rangle$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complement-of-cspan*: $\langle \text{orthogonal-complement } A = \text{orthogonal-complement (cspan } A) \rangle$
 $\langle \text{proof} \rangle$

lemma *orthogonal-complement-orthogonal-complement-closure-cspan*:
 $\langle \text{orthogonal-complement (orthogonal-complement } S) = \text{closure (cspan } S) \rangle$ **for** S
 $:: \langle 'a::\text{hilbert-space set} \rangle$

⟨proof⟩

instance *ccsubspace* :: (*chilbert-space*) *complete-orthomodular-lattice*
⟨proof⟩

9.6 Orthogonal spaces

definition $\langle \text{orthogonal-spaces } S \ T \longleftrightarrow (\forall x \in \text{space-as-set } S. \ \forall y \in \text{space-as-set } T. \text{ is-orthogonal } x \ y) \rangle$

lemma *orthogonal-spaces-leq-compl*: $\langle \text{orthogonal-spaces } S \ T \longleftrightarrow S \leq -T \rangle$
⟨proof⟩

lemma *orthogonal-spaces-bot-right[simp]*: $\langle \text{orthogonal-spaces } S \ \text{bot} \rangle$
⟨proof⟩

lemma *orthogonal-spaces-bot-left[simp]*: $\langle \text{orthogonal-spaces } \text{bot } S \rangle$
⟨proof⟩

lemma *orthogonal-spaces-sym*: $\langle \text{orthogonal-spaces } S \ T \implies \text{orthogonal-spaces } T \ S \rangle$
⟨proof⟩

lemma *orthogonal-sup*: $\langle \text{orthogonal-spaces } S \ T1 \implies \text{orthogonal-spaces } S \ T2 \implies \text{orthogonal-spaces } S \ (\text{sup } T1 \ T2) \rangle$
⟨proof⟩

lemma *orthogonal-sum*:
 assumes $\langle \text{finite } F \rangle$ **and** $\langle \bigwedge x. x \in F \implies \text{orthogonal-spaces } S \ (T \ x) \rangle$
 shows $\langle \text{orthogonal-spaces } S \ (\text{sum } T \ F) \rangle$
 ⟨proof⟩

lemma *orthogonal-spaces-ccspan*: $\langle (\forall x \in S. \ \forall y \in T. \text{ is-orthogonal } x \ y) \longleftrightarrow \text{orthogonal-spaces } (\text{ccspan } S) \ (\text{ccspan } T) \rangle$
⟨proof⟩

lemma *orthogonal-spaces-SUP-left*:
 assumes $\langle \bigwedge x. x \in X \implies \text{orthogonal-spaces } (A \ x) \ B \rangle$
 shows $\langle \text{orthogonal-spaces } (SUP \ x \in X. A \ x) \ B \rangle$
 ⟨proof⟩

lemma *orthogonal-spaces-SUP-right*:
 assumes $\langle \bigwedge x. x \in X \implies \text{orthogonal-spaces } A \ (B \ x) \rangle$
 shows $\langle \text{orthogonal-spaces } A \ (SUP \ x \in X. B \ x) \rangle$
 ⟨proof⟩

9.7 Orthonormal bases

lemma *ortho-basis-exists*:
 fixes $S :: \langle 'a :: \text{chilbert-space set} \rangle$
 assumes $\langle \text{is-ortho-set } S \rangle$

shows $\langle \exists B. B \supseteq S \wedge \text{is-ortho-set } B \wedge \text{closure } (cspan\ B) = UNIV \rangle$
 $\langle \text{proof} \rangle$

lemma *orthonormal-basis-exists*:
fixes $S :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $\langle \text{is-ortho-set } S \rangle$ **and** $\langle \bigwedge x. x \in S \implies \text{norm } x = 1 \rangle$
shows $\langle \exists B. B \supseteq S \wedge \text{is-onb } B \rangle$
 $\langle \text{proof} \rangle$

definition *some-hilbert-basis* :: $\langle 'a::\text{hilbert-space set} \rangle$ **where**
 $\langle \text{some-hilbert-basis} = (SOME\ B::'a\ \text{set. is-onb } B) \rangle$

lemma *is-onb-some-hilbert-basis[simp]*: $\langle \text{is-onb } (\text{some-hilbert-basis} :: 'a::\text{hilbert-space set}) \rangle$
 $\langle \text{proof} \rangle$

lemma *is-ortho-set-some-hilbert-basis[simp]*: $\langle \text{is-ortho-set } \text{some-hilbert-basis} \rangle$
 $\langle \text{proof} \rangle$

lemma *is-normal-some-hilbert-basis*: $\langle \bigwedge x. x \in \text{some-hilbert-basis} \implies \text{norm } x = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *ccspan-some-hilbert-basis[simp]*: $\langle \text{ccspan } \text{some-hilbert-basis} = \text{top} \rangle$
 $\langle \text{proof} \rangle$

lemma *span-some-hilbert-basis[simp]*: $\langle \text{closure } (cspan\ \text{some-hilbert-basis}) = UNIV \rangle$
 $\langle \text{proof} \rangle$

lemma *cindependent-some-hilbert-basis[simp]*: $\langle \text{cindependent } \text{some-hilbert-basis} \rangle$
 $\langle \text{proof} \rangle$

lemma *finite-some-hilbert-basis[simp]*: $\langle \text{finite } (\text{some-hilbert-basis} :: 'a :: \{\text{hilbert-space, cfinite-dim}\}\ \text{set}) \rangle$
 $\langle \text{proof} \rangle$

lemma *some-hilbert-basis-nonempty*: $\langle (\text{some-hilbert-basis} :: 'a::\{\text{hilbert-space, not-singleton}\}\ \text{set}) \neq \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma *basis-projections-reconstruct-has-sum*:
assumes $\langle \text{is-ortho-set } B \rangle$ **and** $\text{norm}B: \langle \bigwedge b. b \in B \implies \text{norm } b = 1 \rangle$ **and** $\psi B: \langle \psi \in \text{space-as-set } (cspan\ B) \rangle$
shows $\langle ((\lambda b. (b \cdot_C \psi) *_C b) \text{ has-sum } \psi)\ B \rangle$
 $\langle \text{proof} \rangle$

lemma *basis-projections-reconstruct*:
assumes $\langle \text{is-ortho-set } B \rangle$ **and** $\langle \bigwedge b. b \in B \implies \text{norm } b = 1 \rangle$ **and** $\langle \psi \in \text{space-as-set } B \rangle$

$\langle \text{ccspan } B \rangle$
shows $\langle (\sum_{b \in B} b \cdot_C \psi) *_C b = \psi \rangle$
 $\langle \text{proof} \rangle$

lemma *basis-projections-reconstruct-summable*:

assumes $\langle \text{is-ortho-set } B \rangle$ **and** $\langle \bigwedge b. b \in B \implies \text{norm } b = 1 \rangle$ **and** $\langle \psi \in \text{space-as-set } (\text{ccspan } B) \rangle$
shows $\langle (\lambda b. (b \cdot_C \psi) *_C b) \text{ summable-on } B \rangle$
 $\langle \text{proof} \rangle$

lemma *parseval-identity-has-sum*:

assumes $\langle \text{is-ortho-set } B \rangle$ **and** $\text{norm} B: \langle \bigwedge b. b \in B \implies \text{norm } b = 1 \rangle$ **and** $\langle \psi \in \text{space-as-set } (\text{ccspan } B) \rangle$
shows $\langle ((\lambda b. (\text{norm } (b \cdot_C \psi))^2) \text{ has-sum } (\text{norm } \psi)^2) B \rangle$
 $\langle \text{proof} \rangle$

lemma *parseval-identity-summable*:

assumes $\langle \text{is-ortho-set } B \rangle$ **and** $\langle \bigwedge b. b \in B \implies \text{norm } b = 1 \rangle$ **and** $\langle \psi \in \text{space-as-set } (\text{ccspan } B) \rangle$
shows $\langle (\lambda b. (\text{norm } (b \cdot_C \psi))^2) \text{ summable-on } B \rangle$
 $\langle \text{proof} \rangle$

lemma *parseval-identity*:

assumes $\langle \text{is-ortho-set } B \rangle$ **and** $\langle \bigwedge b. b \in B \implies \text{norm } b = 1 \rangle$ **and** $\langle \psi \in \text{space-as-set } (\text{ccspan } B) \rangle$
shows $\langle (\sum_{b \in B} (\text{norm } (b \cdot_C \psi))^2) = (\text{norm } \psi)^2 \rangle$
 $\langle \text{proof} \rangle$

9.8 Riesz-representation theorem

lemma *orthogonal-complement-kernel-functional*:

fixes $f :: \langle 'a :: \text{complex-inner} \Rightarrow \text{complex} \rangle$
assumes $\langle \text{bounded-clinear } f \rangle$
shows $\langle \exists x. \text{orthogonal-complement } (f - \cdot \{0\}) = \text{cspan } \{x\} \rangle$
 $\langle \text{proof} \rangle$

lemma *riesz-representation-existence*:

— Theorem 3.4 in [1]
fixes $f :: \langle 'a :: \text{hilbert-space} \Rightarrow \text{complex} \rangle$
assumes $a1: \langle \text{bounded-clinear } f \rangle$
shows $\langle \exists t. \forall x. f x = t \cdot_C x \rangle$
 $\langle \text{proof} \rangle$

lemma *riesz-representation-unique*:

— Theorem 3.4 in [1]
fixes $f :: \langle 'a :: \text{complex-inner} \Rightarrow \text{complex} \rangle$
assumes $\langle \bigwedge x. f x = (t \cdot_C x) \rangle$
assumes $\langle \bigwedge x. f x = (u \cdot_C x) \rangle$
shows $\langle t = u \rangle$

⟨proof⟩

9.9 Adjoints

definition $\langle is-cadjoint\ F\ G \longleftrightarrow (\forall\ x\ y.\ (F\ x \cdot_C\ y) = (x \cdot_C\ G\ y)) \rangle$

lemma *is-adjoint-sym*:

$\langle is-cadjoint\ F\ G \implies is-cadjoint\ G\ F \rangle$

⟨proof⟩

definition $\langle cadjoint\ G = (SOME\ F.\ is-cadjoint\ F\ G) \rangle$
for $G :: 'b::complex-inner \Rightarrow 'a::complex-inner$

lemma *cadjoint-exists*:

fixes $G :: 'b::chilbert-space \Rightarrow 'a::complex-inner$

assumes $[simp]: \langle bounded-clinear\ G \rangle$

shows $\langle \exists\ F.\ is-cadjoint\ F\ G \rangle$

⟨proof⟩

include *norm-syntax*

⟨proof⟩

lemma *cadjoint-is-cadjoint[simp]*:

fixes $G :: 'b::chilbert-space \Rightarrow 'a::complex-inner$

assumes $[simp]: \langle bounded-clinear\ G \rangle$

shows $\langle is-cadjoint\ (cadjoint\ G)\ G \rangle$

⟨proof⟩

lemma *is-cadjoint-unique*:

assumes $\langle is-cadjoint\ F1\ G \rangle$

assumes $\langle is-cadjoint\ F2\ G \rangle$

shows $\langle F1 = F2 \rangle$

⟨proof⟩

lemma *cadjoint-univ-prop*:

fixes $G :: 'b::chilbert-space \Rightarrow 'a::complex-inner$

assumes $a1: \langle bounded-clinear\ G \rangle$

shows $\langle cadjoint\ G\ x \cdot_C\ y = x \cdot_C\ G\ y \rangle$

⟨proof⟩

lemma *cadjoint-univ-prop'*:

fixes $G :: 'b::chilbert-space \Rightarrow 'a::complex-inner$

assumes $a1: \langle bounded-clinear\ G \rangle$

shows $\langle x \cdot_C\ cadjoint\ G\ y = G\ x \cdot_C\ y \rangle$

⟨proof⟩

notation *cadjoint* $(\langle \cdot^\dagger \rangle [99]\ 100)$

lemma *cadjoint-eqI*:

fixes $G :: \langle 'b::complex-inner \Rightarrow 'a::complex-inner \rangle$

```

    and  $F :: \langle 'a \Rightarrow 'b \rangle$ 
    assumes  $\langle \bigwedge x y. (F x \cdot_C y) = (x \cdot_C G y) \rangle$ 
    shows  $\langle G^\dagger = F \rangle$ 
     $\langle proof \rangle$ 

lemma cadjoint-bounded-clinear:
  fixes  $A :: 'a::chilbert-space \Rightarrow 'b::complex-inner$ 
  assumes  $a1: bounded-clinear A$ 
  shows  $\langle bounded-clinear (A^\dagger) \rangle$ 
   $\langle proof \rangle$ 
  include norm-syntax
   $\langle proof \rangle$ 

proposition double-cadjoint:
  fixes  $U :: \langle 'a::chilbert-space \Rightarrow 'b::complex-inner \rangle$ 
  assumes  $a1: bounded-clinear U$ 
  shows  $U^{\dagger\dagger} = U$ 
   $\langle proof \rangle$ 

lemma cadjoint-id[simp]:  $\langle id^\dagger = id \rangle$ 
   $\langle proof \rangle$ 

lemma scaleC-cadjoint:
  fixes  $A :: 'a::chilbert-space \Rightarrow 'b::complex-inner$ 
  assumes  $bounded-clinear A$ 
  shows  $\langle (\lambda t. a *_C A t)^\dagger = (\lambda s. cnj a *_C (A^\dagger) s) \rangle$ 
   $\langle proof \rangle$ 

lemma is-projection-on-is-cadjoint:
  fixes  $M :: \langle 'a::complex-inner set \rangle$ 
  assumes  $a1: \langle is-projection-on \pi M \rangle$  and  $a2: \langle closed-csubspace M \rangle$ 
  shows  $\langle is-cadjoint \pi \pi \rangle$ 
   $\langle proof \rangle$ 

lemma is-projection-on-cadjoint:
  fixes  $M :: \langle 'a::complex-inner set \rangle$ 
  assumes  $\langle is-projection-on \pi M \rangle$  and  $\langle closed-csubspace M \rangle$ 
  shows  $\langle \pi^\dagger = \pi \rangle$ 
   $\langle proof \rangle$ 

lemma projection-cadjoint:
  fixes  $M :: \langle 'a::chilbert-space set \rangle$ 
  assumes  $\langle closed-csubspace M \rangle$ 
  shows  $\langle (projection M)^\dagger = projection M \rangle$ 
   $\langle proof \rangle$ 

```

9.10 More projections

These lemmas logically belong in the "projections" section above but depend on lemmas developed later.

lemma *is-projection-on-plus*:

assumes $\bigwedge x y. x \in A \implies y \in B \implies \text{is-orthogonal } x y$
assumes $\langle \text{closed-csubspace } A \rangle$
assumes $\langle \text{closed-csubspace } B \rangle$
assumes $\langle \text{is-projection-on } \pi A A \rangle$
assumes $\langle \text{is-projection-on } \pi B B \rangle$
shows $\langle \text{is-projection-on } (\lambda x. \pi A x + \pi B x) (A +_M B) \rangle$
 $\langle \text{proof} \rangle$

lemma *projection-plus*:

fixes $A B :: 'a::\text{hilbert-space set}$
assumes $\bigwedge x y. x:A \implies y:B \implies \text{is-orthogonal } x y$
assumes $\langle \text{closed-csubspace } A \rangle$
assumes $\langle \text{closed-csubspace } B \rangle$
shows $\langle \text{projection } (A +_M B) = (\lambda x. \text{projection } A x + \text{projection } B x) \rangle$
 $\langle \text{proof} \rangle$

lemma *is-projection-on-insert*:

assumes *ortho*: $\bigwedge s. s \in S \implies \text{is-orthogonal } a s$
assumes $\langle \text{is-projection-on } \pi (\text{closure } (\text{cspan } S)) \rangle$
assumes $\langle \text{is-projection-on } \pi a (\text{cspan } \{a\}) \rangle$
shows $\text{is-projection-on } (\lambda x. \pi a x + \pi x) (\text{closure } (\text{cspan } (\text{insert } a S)))$
 $\langle \text{proof} \rangle$

lemma *projection-insert*:

fixes $a :: \langle 'a::\text{hilbert-space} \rangle$
assumes $a1: \bigwedge s. s \in S \implies \text{is-orthogonal } a s$
shows $\text{projection } (\text{closure } (\text{cspan } (\text{insert } a S))) u$
 $= \text{projection } (\text{cspan } \{a\}) u + \text{projection } (\text{closure } (\text{cspan } S)) u$
 $\langle \text{proof} \rangle$

lemma *projection-insert-finite*:

fixes $S :: \langle 'a::\text{hilbert-space set} \rangle$
assumes $a1: \bigwedge s. s \in S \implies \text{is-orthogonal } a s$ **and** $a2: \text{finite } S$
shows $\text{projection } (\text{cspan } (\text{insert } a S)) u$
 $= \text{projection } (\text{cspan } \{a\}) u + \text{projection } (\text{cspan } S) u$
 $\langle \text{proof} \rangle$

9.11 Canonical basis (*onb-enum*)

$\langle ML \rangle$

class *onb-enum* = *basis-enum* + *complex-inner* +

assumes *is-orthonormal*: *is-ortho-set* (*set canonical-basis*)
and *is-normal*: $\bigwedge x. x \in (\text{set canonical-basis}) \implies \text{norm } x = 1$

$\langle ML \rangle$

lemma *cinner-canonical-basis*:

assumes $\langle i < \text{length } (\text{canonical-basis} :: 'a::\text{onb-enum list}) \rangle$
assumes $\langle j < \text{length } (\text{canonical-basis} :: 'a::\text{onb-enum list}) \rangle$
shows $\langle \text{cinner } (\text{canonical-basis}!i :: 'a) (\text{canonical-basis}!j) = (\text{if } i=j \text{ then } 1 \text{ else } 0) \rangle$
 $\langle \text{proof} \rangle$

lemma *canonical-basis-is-onb[simp]*: $\langle \text{is-onb } (\text{set } \text{canonical-basis} :: 'a::\text{onb-enum set}) \rangle$
 $\langle \text{proof} \rangle$

instance *onb-enum* $\subseteq$ *hilbert-space*
 $\langle \text{proof} \rangle$

9.12 Conjugate space

instantiation *conjugate-space* :: (*complex-inner*) *complex-inner* **begin**

lift-definition *cinner-conjugate-space* :: $'a \text{ conjugate-space} \Rightarrow 'a \text{ conjugate-space} \Rightarrow \text{complex is}$

$\langle \lambda x y. \text{cinner } y x \rangle \langle \text{proof} \rangle$

instance

$\langle \text{proof} \rangle$

end

instance *conjugate-space* :: (*hilbert-space*) *hilbert-space* $\langle \text{proof} \rangle$

9.13 Misc (ctd.)

lemma *separating-dense-span*:

assumes $\langle \bigwedge F G :: 'a::\text{hilbert-space} \Rightarrow 'b::\{\text{complex-normed-vector}, \text{not-singleton}\}. \text{bounded-clinear } F \Longrightarrow \text{bounded-clinear } G \Longrightarrow (\forall x \in S. F x = G x) \Longrightarrow F = G \rangle$
shows $\langle \text{closure } (\text{cspan } S) = \text{UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *antilinear-cinner*:

shows $\langle \text{antilinear } (\lambda x. x \cdot_C y) \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-extensionality-basis*:

fixes $g h :: 'a::\text{complex-inner}$
assumes $\langle \text{ccspan } B = \text{top} \rangle$
assumes $\langle \bigwedge x. x \in B \Longrightarrow x \cdot_C g = x \cdot_C h \rangle$
shows $\langle g = h \rangle$
 $\langle \text{proof} \rangle$

end

10 *One-Dimensional-Spaces* – One dimensional complex vector spaces

theory *One-Dimensional-Spaces*

imports

Complex-Inner-Product

Complex-Bounded-Operators.Extra-Operator-Norm

begin

The class *one-dim* applies to one-dimensional vector spaces. Those are additionally interpreted as *complex-algebra-1*s via the canonical isomorphism between a one-dimensional vector space and *complex*.

class *one-dim* = *onb-enum* + *one* + *times* + *inverse* +
assumes *one-dim-canonical-basis[simp]*: *canonical-basis* = [1]
assumes *one-dim-prod-scale1*: $(a *_C 1) * (b *_C 1) = (a * b) *_C 1$
assumes *divide-inverse*: $x / y = x * \text{inverse } y$
assumes *one-dim-inverse*: $\text{inverse } (a *_C 1) = \text{inverse } a *_C 1$

hide-fact (open) *divide-inverse*

— *divide-inverse* from class *field*, instantiated below, subsumes this fact.

instance *complex* :: *one-dim*

⟨*proof*⟩

lemma *one-cinner-one[simp]*: $\langle (1 :: 'a :: \text{one-dim}) \cdot_C 1 = 1 \rangle$

⟨*proof*⟩

include *norm-syntax*

⟨*proof*⟩

lemma *one-cinner-a-scaleC-one[simp]*: $\langle ((1 :: 'a :: \text{one-dim}) \cdot_C a) *_C 1 = a \rangle$

⟨*proof*⟩

lemma *one-dim-apply-is-times-def*:

$\psi * \varphi = ((1 \cdot_C \psi) * (1 \cdot_C \varphi)) *_C 1$ **for** $\psi :: 'a :: \text{one-dim}$

⟨*proof*⟩

instance *one-dim* $\subseteq$ *complex-algebra-1*

⟨*proof*⟩

instance *one-dim* $\subseteq$ *complex-normed-algebra*

⟨*proof*⟩

instance *one-dim* $\subseteq$ *complex-normed-algebra-1*

⟨*proof*⟩

This is the canonical isomorphism between any two one dimensional spaces. Specifically, if 1 denotes the element of the canonical basis (which is specified by type class *basis-enum*), then *one-dim-iso* is the unique isomorphism that maps 1 to 1 .

definition *one-dim-iso* :: 'a::one-dim $\Rightarrow$ 'b::one-dim
where *one-dim-iso* a = of-complex (1 $\cdot_C$ a)

lemma *one-dim-iso-idem*[simp]: *one-dim-iso* (*one-dim-iso* x) = *one-dim-iso* x
 ⟨proof⟩

lemma *one-dim-iso-id*[simp]: *one-dim-iso* = id
 ⟨proof⟩

lemma *one-dim-iso-adjoint*[simp]: $\langle \text{cadjoint } \textit{one-dim-iso} = \textit{one-dim-iso} \rangle$
 ⟨proof⟩

lemma *one-dim-iso-is-of-complex*[simp]: *one-dim-iso* = of-complex
 ⟨proof⟩

lemma *of-complex-one-dim-iso*[simp]: of-complex (*one-dim-iso* ψ) = *one-dim-iso* ψ
 ⟨proof⟩

lemma *one-dim-iso-of-complex*[simp]: *one-dim-iso* (of-complex c) = of-complex c
 ⟨proof⟩

lemma *one-dim-iso-add*[simp]:
 $\langle \textit{one-dim-iso} (a + b) = \textit{one-dim-iso} a + \textit{one-dim-iso} b \rangle$
 ⟨proof⟩

lemma *one-dim-iso-minus*[simp]:
 $\langle \textit{one-dim-iso} (a - b) = \textit{one-dim-iso} a - \textit{one-dim-iso} b \rangle$
 ⟨proof⟩

lemma *one-dim-iso-scaleC*[simp]: *one-dim-iso* (c $\cdot_C$ ψ) = c $\cdot_C$ *one-dim-iso* ψ
 ⟨proof⟩

lemma *clinear-one-dim-iso*[simp]: *clinear* *one-dim-iso*
 ⟨proof⟩

lemma *bounded-clinear-one-dim-iso*[simp]: *bounded-clinear* *one-dim-iso*
 ⟨proof⟩

lemma *one-dim-iso-of-one*[simp]: *one-dim-iso* 1 = 1
 ⟨proof⟩

lemma *onorm-one-dim-iso*[simp]: *onorm* *one-dim-iso* = 1
 ⟨proof⟩

lemma *one-dim-iso-times[simp]*: *one-dim-iso* $(\psi * \varphi) = \text{one-dim-iso } \psi * \text{one-dim-iso } \varphi$

$\langle \text{proof} \rangle$

lemma *one-dim-iso-of-zero[simp]*: *one-dim-iso* $0 = 0$

$\langle \text{proof} \rangle$

lemma *one-dim-iso-of-zero'*: *one-dim-iso* $x = 0 \implies x = 0$

$\langle \text{proof} \rangle$

lemma *one-dim-scaleC-1[simp]*: *one-dim-iso* $x *_C 1 = x$

$\langle \text{proof} \rangle$

lemma *one-dim-clinear-eqI*:

assumes $(x::'a::\text{one-dim}) \neq 0$ **and** *clinear* f **and** *clinear* g **and** $f\ x = g\ x$

shows $f = g$

$\langle \text{proof} \rangle$

lemma *one-dim-norm*: *norm* $x = \text{cmod } (\text{one-dim-iso } x)$

$\langle \text{proof} \rangle$

lemma *norm-one-dim-iso[simp]*: $\langle \text{norm } (\text{one-dim-iso } x) = \text{norm } x \rangle$

$\langle \text{proof} \rangle$

lemma *one-dim-onorm*:

fixes $f :: 'a::\text{one-dim} \Rightarrow 'b::\text{complex-normed-vector}$

assumes *clinear* f

shows *onorm* $f = \text{norm } (f\ 1)$

$\langle \text{proof} \rangle$

lemma *one-dim-onorm'*:

fixes $f :: 'a::\text{one-dim} \Rightarrow 'b::\text{one-dim}$

assumes *clinear* f

shows *onorm* $f = \text{cmod } (\text{one-dim-iso } (f\ 1))$

$\langle \text{proof} \rangle$

instance *one-dim* $\subseteq$ *zero-neq-one* $\langle \text{proof} \rangle$

lemma *one-dim-iso-inj*: *one-dim-iso* $x = \text{one-dim-iso } y \implies x = y$

$\langle \text{proof} \rangle$

instance *one-dim* $\subseteq$ *comm-ring*

$\langle \text{proof} \rangle$

instance *one-dim* $\subseteq$ *field*

$\langle \text{proof} \rangle$

instance *one-dim* $\subseteq$ *complex-normed-field*

$\langle \text{proof} \rangle$

instance *one-dim* $\subseteq$ *hilbert-space* $\langle \text{proof} \rangle$

lemma *ccspan-one-dim[simp]*: $\langle \text{ccspan } \{x\} = \text{top} \rangle$ **if** $\langle x \neq 0 \rangle$ **for** $x :: \langle - :: \text{one-dim} \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-ccsubspace-all-or-nothing*: $\langle A = \text{bot} \vee A = \text{top} \rangle$ **for** $A :: \langle - :: \text{one-dim ccsubspace} \rangle$
 $\langle \text{proof} \rangle$

lemma *scaleC-1-right[simp]*: $\langle \text{scaleC } x (1 :: 'a :: \text{one-dim}) = \text{of-complex } x \rangle$
 $\langle \text{proof} \rangle$

lemma *canonical-basis-length-one-dim[simp]*: $\langle \text{canonical-basis-length TYPE('a :: \text{one-dim})} = 1 \rangle$
 $\langle \text{proof} \rangle$

end

11 Complex-Euclidean-Space0 – Finite-Dimensional Inner Product Spaces

theory *Complex-Euclidean-Space0*

imports

HOL-Analysis.L2-Norm

Complex-Inner-Product

HOL-Analysis.Product-Vector

HOL-Library.Rewrite

begin

11.1 Type class of Euclidean spaces

class *euclidean-space* = *complex-inner* +

fixes *CBasis* :: 'a set

assumes *nonempty-CBasis* [simp]: *CBasis* $\neq \{\}$

assumes *finite-CBasis* [simp]: *finite CBasis*

assumes *cinner-CBasis*:

$\llbracket u \in \text{CBasis}; v \in \text{CBasis} \rrbracket \implies \text{cinner } u \ v = (\text{if } u = v \text{ then } 1 \text{ else } 0)$

assumes *euclidean-all-zero-iff*:

$(\forall u \in \text{CBasis}. \text{cinner } x \ u = 0) \longleftrightarrow (x = 0)$

syntax *-type-cdimension* :: *type* $\Rightarrow$ *nat* $(\langle (1 \text{CDIM} / (1'(-))) \rangle)$

syntax-consts *-type-cdimension* $\equiv$ *card*

translations *CDIM*('a) $\rightarrow$ *CONST card* (*CONST CBasis* :: 'a set)

$\langle \text{ML} \rangle$

lemma (**in** *euclidean-space*) *norm-CBasis*[simp]: $u \in \text{CBasis} \implies \text{norm } u = 1$

$\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *cinner-same-CBasis*[simp]: $u \in \text{CBasis} \implies \text{cinner } u \ u = 1$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *cinner-not-same-CBasis*: $u \in \text{CBasis} \implies v \in \text{CBasis} \implies u \neq v \implies \text{cinner } u \ v = 0$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *sgn-CBasis*: $u \in \text{CBasis} \implies \text{sgn } u = u$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *CBasis-zero* [simp]: $0 \notin \text{CBasis}$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *nonzero-CBasis*: $u \in \text{CBasis} \implies u \neq 0$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *SOME-CBasis*: $(\text{SOME } i. i \in \text{CBasis}) \in \text{CBasis}$
 $\langle \text{proof} \rangle$

lemma *norm-some-CBasis* [simp]: $\text{norm } (\text{SOME } i. i \in \text{CBasis}) = 1$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *cinner-sum-left-CBasis*[simp]:
 $b \in \text{CBasis} \implies \text{cinner } (\sum i \in \text{CBasis}. f \ i *_{\text{C}} \ i) \ b = \text{cnj } (f \ b)$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *ceukclidean-eqI*:
assumes $b: \bigwedge b. b \in \text{CBasis} \implies \text{cinner } x \ b = \text{cinner } y \ b$ **shows** $x = y$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *ceukclidean-eq-iff*:
 $x = y \longleftrightarrow (\forall b \in \text{CBasis}. \text{cinner } x \ b = \text{cinner } y \ b)$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *ceukclidean-representation-sum*:
 $(\sum i \in \text{CBasis}. f \ i *_{\text{C}} \ i) = b \longleftrightarrow (\forall i \in \text{CBasis}. f \ i = \text{cnj } (\text{cinner } b \ i))$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *ceukclidean-representation-sum'*:
 $b = (\sum i \in \text{CBasis}. f \ i *_{\text{C}} \ i) \longleftrightarrow (\forall i \in \text{CBasis}. f \ i = \text{cinner } i \ b)$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *ceukclidean-representation*: $(\sum b \in \text{CBasis}. \text{cinner } b \ x$

$*_C b) = x$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *ceukclidean-cinner*: $\text{cinner } x \ y = (\sum_{b \in \text{CBasis}} \text{cinner } x \ b * \text{cnj } (\text{cinner } y \ b))$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *choice-CBasis-iff*:
fixes $P :: 'a \Rightarrow \text{complex} \Rightarrow \text{bool}$
shows $(\forall i \in \text{CBasis}. \exists x. P \ i \ x) \longleftrightarrow (\exists x. \forall i \in \text{CBasis}. P \ i \ (\text{cinner } x \ i))$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *bchoice-CBasis-iff*:
fixes $P :: 'a \Rightarrow \text{complex} \Rightarrow \text{bool}$
shows $(\forall i \in \text{CBasis}. \exists x \in A. P \ i \ x) \longleftrightarrow (\exists x. \forall i \in \text{CBasis}. \text{cinner } x \ i \in A \wedge P \ i \ (\text{cinner } x \ i))$
 $\langle \text{proof} \rangle$

lemma (in *ceukclidean-space*) *ceukclidean-representation-sum-fun*:
 $(\lambda x. \sum_{b \in \text{CBasis}} \text{cinner } b \ (f \ x) *_C b) = f$
 $\langle \text{proof} \rangle$

lemma *euclidean-isCont*:
assumes $\bigwedge b. b \in \text{CBasis} \implies \text{isCont } (\lambda x. (\text{cinner } b \ (f \ x)) *_C b) \ x$
shows $\text{isCont } f \ x$
 $\langle \text{proof} \rangle$

lemma *CDIM-positive* [simp]: $0 < \text{CDIM}('a::\text{ceukclidean-space})$
 $\langle \text{proof} \rangle$

lemma *CDIM-ge-Suc0* [simp]: $\text{Suc } 0 \leq \text{card } \text{CBasis}$
 $\langle \text{proof} \rangle$

lemma *sum-cinner-CBasis-scaleC* [simp]:
fixes $f :: 'a::\text{ceukclidean-space} \Rightarrow 'b::\text{complex-vector}$
assumes $b \in \text{CBasis}$ **shows** $(\sum_{i \in \text{CBasis}} (\text{cinner } i \ b) *_C f \ i) = f \ b$
 $\langle \text{proof} \rangle$

lemma *sum-cinner-CBasis-eq* [simp]:
assumes $b \in \text{CBasis}$ **shows** $(\sum_{i \in \text{CBasis}} (\text{cinner } i \ b) * f \ i) = f \ b$
 $\langle \text{proof} \rangle$

lemma *sum-if-cinner* [simp]:
assumes $i \in \text{CBasis} \ j \in \text{CBasis}$
shows $\text{cinner } (\sum_{k \in \text{CBasis}} \text{if } k = i \text{ then } f \ i *_C i \text{ else } g \ k *_C k) \ j = (\text{if } j = i \text{ then } \text{cnj } (f \ j) \text{ else } \text{cnj } (g \ j))$
 $\langle \text{proof} \rangle$

lemma *norm-le-componentwise*:

$(\bigwedge b. b \in CBasis \implies cmod(cinner\ x\ b) \leq cmod(cinner\ y\ b)) \implies norm\ x \leq norm\ y$
 $\langle proof \rangle$

lemma *CBasis-le-norm*: $b \in CBasis \implies cmod\ (cinner\ x\ b) \leq norm\ x$
 $\langle proof \rangle$

lemma *norm-bound-CBasis-le*: $b \in CBasis \implies norm\ x \leq e \implies cmod\ (inner\ x\ b) \leq e$
 $\langle proof \rangle$

lemma *norm-bound-CBasis-lt*: $b \in CBasis \implies norm\ x < e \implies cmod\ (inner\ x\ b) < e$
 $\langle proof \rangle$

lemma *cnorm-le-l1*: $norm\ x \leq (\sum b \in CBasis. cmod\ (cinner\ x\ b))$
 $\langle proof \rangle$

11.2 Class instances

11.2.1 Type *complex*

instantiation *complex* :: *ceclidean-space*
begin

definition

$[simp]: CBasis = \{1::complex\}$

instance

$\langle proof \rangle$

end

lemma *CDIM-complex* $[simp]$: $CDIM(complex) = 1$
 $\langle proof \rangle$

11.2.2 Type $'a \times 'b$

lemma *cinner-Pair* $[simp]$: $cinner\ (a, b)\ (c, d) = cinner\ a\ c + cinner\ b\ d$
 $\langle proof \rangle$

lemma *cinner-Pair-0*: $cinner\ x\ (0, b) = cinner\ (snd\ x)\ b$ $cinner\ x\ (a, 0) = cinner\ (fst\ x)\ a$
 $\langle proof \rangle$

instantiation *prod* :: (*ceclidean-space*, *ceclidean-space*) *ceclidean-space*
begin

definition

```

CBasis = (λu. (u, 0)) ‘ CBasis ∪ (λv. (0, v)) ‘ CBasis

lemma sum-CBasis-prod-eq:
  fixes f::('a*'b)⇒('a*'b)
  shows sum f CBasis = sum (λi. f (i, 0)) CBasis + sum (λi. f (0, i)) CBasis
  ⟨proof⟩

instance ⟨proof⟩

lemma CDIM-prod[simp]: CDIM('a × 'b) = CDIM('a) + CDIM('b)
  ⟨proof⟩

end

```

11.3 Locale instances

```

lemma finite-dimensional-vector-space-euclidean:
  finite-dimensional-vector-space (*C) CBasis
  ⟨proof⟩

interpretation ceubl: finite-dimensional-vector-space scaleC :: complex => 'a =>
'a::euclidean-space CBasis
  rewrites module.dependent (*C) = cdependent
    and module.representation (*C) = crepresentation
    and module.subspace (*C) = csubspace
    and module.span (*C) = cspan
    and vector-space.extend-basis (*C) = certend-basis
    and vector-space.dim (*C) = cdim
    and Vector-Spaces.linear (*C) (*C) = clinear
    and Vector-Spaces.linear (*) (*C) = clinear
    and finite-dimensional-vector-space.dimension CBasis = CDIM('a)

  ⟨proof⟩

interpretation ceubl: finite-dimensional-vector-space-pair-1
  scaleC::complex⇒'a::euclidean-space⇒'a CBasis
  scaleC::complex⇒'b::complex-vector ⇒ 'b
  ⟨proof⟩

interpretation ceubl?: finite-dimensional-vector-space-prod scaleC scaleC CBasis
CBasis
  rewrites Basis-pair = CBasis
    and module-prod.scale (*C) (*C) = (scaleC::=>=>('a × 'b))
  ⟨proof⟩

end

```


12 Complex-Bounded-Linear-Function0 – Bounded Linear Function

```

theory Complex-Bounded-Linear-Function0
imports
  HOL-Analysis.Bounded-Linear-Function
  Complex-Inner-Product
  Complex-Euclidean-Space0
begin

```

```

unbundle cinner-syntax

```

```

lemma conorm-componentwise:
  assumes bounded-clinear  $f$ 
  shows  $\text{onorm } f \leq (\sum i \in \text{CBasis}. \text{norm } (f\ i))$ 
   $\langle \text{proof} \rangle$ 

```

```

lemmas conorm-componentwise-le = order-trans[OF conorm-componentwise]

```

12.1 Intro rules for bounded-linear

```

lemma onorm-cinner-left:
  assumes bounded-linear  $r$ 
  shows  $\text{onorm } (\lambda x. r\ x \cdot_C f) \leq \text{onorm } r * \text{norm } f$ 
   $\langle \text{proof} \rangle$ 

```

```

lemma onorm-cinner-right:
  assumes bounded-linear  $r$ 
  shows  $\text{onorm } (\lambda x. f \cdot_C r\ x) \leq \text{norm } f * \text{onorm } r$ 
   $\langle \text{proof} \rangle$ 

```

```

lemmas [bounded-linear-intros] =
  bounded-clinear-zero
  bounded-clinear-add
  bounded-clinear-const-mult
  bounded-clinear-mult-const
  bounded-clinear-scaleC-const
  bounded-clinear-const-scaleC
  bounded-clinear-const-scaleR
  bounded-clinear-ident
  bounded-clinear-sum

```

```

  bounded-clinear-sub

```

```

  bounded-antilinear-cinner-left-comp
  bounded-clinear-cinner-right-comp

```

12.2 declaration of derivative/continuous/tendsto introduction rules for bounded linear functions

$\langle ML \rangle$

12.3 Type of complex bounded linear functions

typedef (overloaded) ('a, 'b) cblinfun ($\langle (- \Rightarrow_{CL} /-) \rangle$ [22, 21] 21) =
 {f::'a::complex-normed-vector $\Rightarrow$ 'b::complex-normed-vector. bounded-clinear f}
morphisms cblinfun-apply CBlinfun
 $\langle proof \rangle$

declare [[coercion
 cblinfun-apply :: ('a::complex-normed-vector $\Rightarrow_{CL}$ 'b::complex-normed-vector)
 $\Rightarrow$ 'a $\Rightarrow$ 'b]]

lemma bounded-clinear-cblinfun-apply[bounded-linear-intros]:
 bounded-clinear g $\implies$ bounded-clinear ($\lambda x. cblinfun-apply f (g x)$)
 $\langle proof \rangle$

setup-lifting type-definition-cblinfun

lemma cblinfun-eqI: ($\bigwedge i. cblinfun-apply x i = cblinfun-apply y i$) $\implies x = y$
 $\langle proof \rangle$

lemma bounded-clinear-CBlinfun-apply: bounded-clinear f $\implies cblinfun-apply (CBlinfun f) = f$
 $\langle proof \rangle$

12.4 Type class instantiations

instantiation cblinfun :: (complex-normed-vector, complex-normed-vector) complex-normed-vector
begin

lift-definition norm-cblinfun :: 'a $\Rightarrow_{CL}$ 'b $\Rightarrow$ real **is** onorm $\langle proof \rangle$

lift-definition minus-cblinfun :: 'a $\Rightarrow_{CL}$ 'b $\Rightarrow$ 'a $\Rightarrow_{CL}$ 'b $\Rightarrow$ 'a $\Rightarrow_{CL}$ 'b
is $\lambda f g x. f x - g x$
 $\langle proof \rangle$

definition dist-cblinfun :: 'a $\Rightarrow_{CL}$ 'b $\Rightarrow$ 'a $\Rightarrow_{CL}$ 'b $\Rightarrow$ real
where dist-cblinfun a b = norm (a - b)

definition [code del]:
 (uniformity :: (('a $\Rightarrow_{CL}$ 'b) $\times$ ('a $\Rightarrow_{CL}$ 'b)) filter) = (INF e $\in$ {0 <..}. principal { (x, y). dist x y < e })

definition open-cblinfun :: ('a $\Rightarrow_{CL}$ 'b) set $\Rightarrow$ bool

where $[code\ del]: open-cblinfun\ S = (\forall x \in S. \forall_F (x', y) \text{ in uniformity. } x' = x \longrightarrow y \in S)$

lift-definition $uminus-cblinfun :: 'a \Rightarrow_{CL} 'b \Rightarrow 'a \Rightarrow_{CL} 'b$ **is** $\lambda f\ x. - f\ x$
 $\langle proof \rangle$

lift-definition $zero-cblinfun :: 'a \Rightarrow_{CL} 'b$ **is** $\lambda x. 0$
 $\langle proof \rangle$

lift-definition $plus-cblinfun :: 'a \Rightarrow_{CL} 'b \Rightarrow 'a \Rightarrow_{CL} 'b \Rightarrow 'a \Rightarrow_{CL} 'b$
is $\lambda f\ g\ x. f\ x + g\ x$
 $\langle proof \rangle$

lift-definition $scaleC-cblinfun::complex \Rightarrow 'a \Rightarrow_{CL} 'b \Rightarrow 'a \Rightarrow_{CL} 'b$ **is** $\lambda r\ f\ x. r$
 $*_C\ f\ x$
 $\langle proof \rangle$

lift-definition $scaleR-cblinfun::real \Rightarrow 'a \Rightarrow_{CL} 'b \Rightarrow 'a \Rightarrow_{CL} 'b$ **is** $\lambda r\ f\ x. r *_R f\ x$
 $\langle proof \rangle$

definition $sgn-cblinfun :: 'a \Rightarrow_{CL} 'b \Rightarrow 'a \Rightarrow_{CL} 'b$
where $sgn-cblinfun\ x = scaleC\ (inverse\ (norm\ x))\ x$

instance
 $\langle proof \rangle$

end

declare $uniformity-Absort$ [**where** $'a = ('a :: complex-normed-vector) \Rightarrow_{CL} ('b :: complex-normed-vector),\ code]$

lemma $norm-cblinfun-eqI$:
assumes $n \leq norm\ (cblinfun-apply\ f\ x) / norm\ x$
assumes $\bigwedge x. norm\ (cblinfun-apply\ f\ x) \leq n * norm\ x$
assumes $0 \leq n$
shows $norm\ f = n$
 $\langle proof \rangle$

lemma $norm-cblinfun$: $norm\ (cblinfun-apply\ f\ x) \leq norm\ f * norm\ x$
 $\langle proof \rangle$

lemma $norm-cblinfun-bound$: $0 \leq b \implies (\bigwedge x. norm\ (cblinfun-apply\ f\ x) \leq b * norm\ x) \implies norm\ f \leq b$
 $\langle proof \rangle$

lemma $bounded-cbilinear-cblinfun-apply[bounded-cbilinear]$: $bounded-cbilinear\ cblinfun-apply$
 $\langle proof \rangle$

interpretation $cblinfun$: $bounded-cbilinear\ cblinfun-apply$

```

    <proof>

lemmas bounded-clinear-apply-cblinfun [intro, simp] = cblinfun.bounded-clinear-left

declare cblinfun.zero-left [simp] cblinfun.zero-right [simp]

context bounded-cbilinear
begin

named-theorems cbilinear-simps

lemmas [cbilinear-simps] =
  add-left
  add-right
  diff-left
  diff-right
  minus-left
  minus-right
  scaleC-left
  scaleC-right
  zero-left
  zero-right
  sum-left
  sum-right

end

instance cblinfun :: (complex-normed-vector, cbanach) cbanach

    <proof>

```

12.5 On Euclidean Space

```

lemma norm-cblinfun-ceuclidean-le:
  fixes a::'a::ceuclidean-space  $\Rightarrow_{CL}$  'b::complex-normed-vector
  shows  $\text{norm } a \leq \text{sum } (\lambda x. \text{norm } (a \ x)) \text{ } CBasis$ 
    <proof>

lemma ctendsto-componentwise1:
  fixes a::'a::ceuclidean-space  $\Rightarrow_{CL}$  'b::complex-normed-vector
  and b::'c  $\Rightarrow$  'a  $\Rightarrow_{CL}$  'b
  assumes  $(\bigwedge j. j \in CBasis \implies ((\lambda n. b \ n \ j) \longrightarrow a \ j) \ F)$ 
  shows  $(b \longrightarrow a) \ F$ 
    <proof>

lift-definition
  cblinfun-of-matrix::('b::ceuclidean-space  $\Rightarrow$  'a::ceuclidean-space  $\Rightarrow$  complex)  $\Rightarrow$  'a

```

$\Rightarrow_{CL} 'b$
is $\lambda a\ x. \sum_{i \in CBasis}. \sum_{j \in CBasis}. ((j \cdot_C x) * a\ i\ j) *_C i$
 $\langle proof \rangle$

lemma *cblinfun-of-matrix-works*:
fixes $f::'a::\text{euclidean-space} \Rightarrow_{CL} 'b::\text{euclidean-space}$
shows $cblinfun\text{-of-matrix}\ (\lambda i\ j. i \cdot_C (f\ j)) = f$
 $\langle proof \rangle$

lemma *cblinfun-of-matrix-apply*:
 $cblinfun\text{-of-matrix}\ a\ x = (\sum_{i \in CBasis}. \sum_{j \in CBasis}. ((j \cdot_C x) * a\ i\ j) *_C i)$
 $\langle proof \rangle$

lemma *cblinfun-of-matrix-minus*: $cblinfun\text{-of-matrix}\ x - cblinfun\text{-of-matrix}\ y =$
 $cblinfun\text{-of-matrix}\ (x - y)$
 $\langle proof \rangle$

lemma *norm-cblinfun-of-matrix*:
 $norm\ (cblinfun\text{-of-matrix}\ a) \leq (\sum_{i \in CBasis}. \sum_{j \in CBasis}. cmod\ (a\ i\ j))$
 $\langle proof \rangle$

lemma *tendsto-cblinfun-of-matrix*:
assumes $\bigwedge i\ j. i \in CBasis \implies j \in CBasis \implies ((\lambda n. b\ n\ i\ j) \longrightarrow a\ i\ j)\ F$
shows $((\lambda n. cblinfun\text{-of-matrix}\ (b\ n)) \longrightarrow cblinfun\text{-of-matrix}\ a)\ F$
 $\langle proof \rangle$

lemma *ctendsto-componentwise*:
fixes $a::'a::\text{euclidean-space} \Rightarrow_{CL} 'b::\text{euclidean-space}$
and $b::'c \Rightarrow 'a \Rightarrow_{CL} 'b$
shows $(\bigwedge i\ j. i \in CBasis \implies j \in CBasis \implies ((\lambda n. b\ n\ j \cdot_C i) \longrightarrow a\ j \cdot_C i)$
 $F) \implies (b \longrightarrow a)\ F$
 $\langle proof \rangle$

lemma
continuous-cblinfun-componentwiseI:
fixes $f::'b::t2\text{-space} \Rightarrow 'a::\text{euclidean-space} \Rightarrow_{CL} 'c::\text{euclidean-space}$
assumes $\bigwedge i\ j. i \in CBasis \implies j \in CBasis \implies \text{continuous}\ F\ (\lambda x. (f\ x)\ j \cdot_C i)$
shows $\text{continuous}\ F\ f$
 $\langle proof \rangle$

lemma
continuous-cblinfun-componentwiseI1:
fixes $f::'b::t2\text{-space} \Rightarrow 'a::\text{euclidean-space} \Rightarrow_{CL} 'c::\text{complex-normed-vector}$
assumes $\bigwedge i. i \in CBasis \implies \text{continuous}\ F\ (\lambda x. f\ x\ i)$
shows $\text{continuous}\ F\ f$
 $\langle proof \rangle$

lemma

continuous-on-cblinfun-componentwise:

fixes $f:: 'd::t2\text{-space} \Rightarrow 'e::\text{euclidean-space} \Rightarrow_{CL} 'f::\text{complex-normed-vector}$

assumes $\bigwedge i. i \in CBasis \implies \text{continuous-on } s \ (\lambda x. f \ x \ i)$

shows *continuous-on* $s \ f$

<proof>

lemma *bounded-antilinear-cblinfun-matrix:* *bounded-antilinear* $(\lambda x. (x::\Rightarrow_{CL} -) \ j \cdot_C i)$

<proof>

lemma *continuous-cblinfun-matrix:*

fixes $f:: 'b::t2\text{-space} \Rightarrow 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'c::\text{complex-inner}$

assumes *continuous* $F \ f$

shows *continuous* $F \ (\lambda x. (f \ x) \ j \cdot_C i)$

<proof>

lemma *continuous-on-cblinfun-matrix:*

fixes $f:: 'a::t2\text{-space} \Rightarrow 'b::\text{complex-normed-vector} \Rightarrow_{CL} 'c::\text{complex-inner}$

assumes *continuous-on* $S \ f$

shows *continuous-on* $S \ (\lambda x. (f \ x) \ j \cdot_C i)$

<proof>

lemma *continuous-on-cblinfun-of-matrix*[*continuous-intros*]:

assumes $\bigwedge i \ j. i \in CBasis \implies j \in CBasis \implies \text{continuous-on } S \ (\lambda s. g \ s \ i \ j)$

shows *continuous-on* $S \ (\lambda s. \text{cblinfun-of-matrix } (g \ s))$

<proof>

lemma *cblinfun-euclidean-eqI:* $(\bigwedge i. i \in CBasis \implies \text{cblinfun-apply } x \ i = \text{cblinfun-apply } y \ i) \implies x = y$

<proof>

lemma *CBlinfun-eq-matrix:* *bounded-clinear* $f \implies CBlinfun \ f = \text{cblinfun-of-matrix } (\lambda i \ j. i \cdot_C f \ j)$

<proof>

12.6 concrete bounded linear functions

lemma *transfer-bounded-cbilinear-bounded-clinearI:*

assumes $g = (\lambda i \ x. (\text{cblinfun-apply } (f \ i) \ x))$

shows *bounded-cbilinear* $g = \text{bounded-clinear } f$

<proof>

lemma *transfer-bounded-cbilinear-bounded-clinear*[*transfer-rule*]:
 (rel-fun (rel-fun (=) (pcr-cblinfun (=) (=))) (=)) bounded-cbilinear bounded-clinear
 ⟨proof⟩

lemma *transfer-bounded-sesquilinear-bounded-antilinearI*:
 assumes $g = (\lambda i\ x. \text{cblinfun-apply } (f\ i)\ x)$
 shows *bounded-sesquilinear* $g = \text{bounded-antilinear } f$
 ⟨proof⟩

lemma *transfer-bounded-sesquilinear-bounded-antilinear*[*transfer-rule*]:
 (rel-fun (rel-fun (=) (pcr-cblinfun (=) (=))) (=)) bounded-sesquilinear bounded-antilinear
 ⟨proof⟩

context *bounded-cbilinear*
begin

lift-definition *prod-left*:: $'b \Rightarrow 'a \Rightarrow_{CL} 'c$ **is** $(\lambda b\ a. \text{prod } a\ b)$
 ⟨proof⟩
declare *prod-left.rep-eq*[*simp*]

lemma *bounded-clinear-prod-left*[*bounded-clinear*]: *bounded-clinear prod-left*
 ⟨proof⟩

lift-definition *prod-right*:: $'a \Rightarrow 'b \Rightarrow_{CL} 'c$ **is** $(\lambda a\ b. \text{prod } a\ b)$
 ⟨proof⟩
declare *prod-right.rep-eq*[*simp*]

lemma *bounded-clinear-prod-right*[*bounded-clinear*]: *bounded-clinear prod-right*
 ⟨proof⟩

end

lift-definition *id-cblinfun*:: $'a::\text{complex-normed-vector} \Rightarrow_{CL} 'a$ **is** $\lambda x. x$
 ⟨proof⟩

lemmas *cblinfun-id-cblinfun-apply*[*simp*] = *id-cblinfun.rep-eq*

lemma *norm-cblinfun-id*[*simp*]:
 $\text{norm } (id-cblinfun::'a::\{\text{complex-normed-vector}, \text{not-singleton}\} \Rightarrow_{CL} 'a) = 1$
 ⟨proof⟩

lemma *norm-cblinfun-id-le*:
 $\text{norm } (id-cblinfun::'a::\text{complex-normed-vector} \Rightarrow_{CL} 'a) \leq 1$
 ⟨proof⟩

lift-definition *cblinfun-compose*::

'a::complex-normed-vector $\Rightarrow_{CL}$ 'b::complex-normed-vector $\Rightarrow$
'c::complex-normed-vector $\Rightarrow_{CL}$ 'a $\Rightarrow$
'c $\Rightarrow_{CL}$ 'b (infixl 67) is (o)

parametric *comp-transfer*

<proof>

lemma *cblinfun-apply-cblinfun-compose[simp]*: $(a \circ_{CL} b) \ c = a \ (b \ c)$

<proof>

lemma *norm-cblinfun-compose*:

$\text{norm } (f \circ_{CL} g) \leq \text{norm } f * \text{norm } g$

<proof>

lemma *bounded-cbilinear-cblinfun-compose[bounded-cbilinear]*: *bounded-cbilinear* $(\circ_{CL})$

<proof>

lemma *cblinfun-compose-zero[simp]*:

blinfun-compose 0 = (λ -. 0)

blinfun-compose x 0 = 0

<proof>

lemma *cblinfun-bij2*:

fixes *f::'a $\Rightarrow_{CL}$ 'a::ceclidean-space*

assumes *f $\circ_{CL}$ g = id-cblinfun*

shows *bij (cblinfun-apply g)*

<proof>

lemma *cblinfun-bij1*:

fixes *f::'a $\Rightarrow_{CL}$ 'a::ceclidean-space*

assumes *f $\circ_{CL}$ g = id-cblinfun*

shows *bij (cblinfun-apply f)*

<proof>

lift-definition *cblinfun-cinner-right::'a::complex-inner $\Rightarrow$ 'a $\Rightarrow_{CL}$ complex* **is** $(\cdot_C)$

<proof>

declare *cblinfun-cinner-right.rep-eq[simp]*

lemma *bounded-antilinear-cblinfun-cinner-right[bounded-antilinear]*: *bounded-antilinear* *cblinfun-cinner-right*

$\langle \text{proof} \rangle$

lift-definition $\text{cblinfun-scale } C\text{-right} :: \text{complex} \Rightarrow 'a \Rightarrow_{CL} 'a :: \text{complex-normed-vector}$
is $(*_C)$
 $\langle \text{proof} \rangle$
declare $\text{cblinfun-scale } C\text{-right.rep-eq}[simp]$

lemma $\text{bounded-clinear-cblinfun-scale } C\text{-right}[\text{bounded-clinear}]$: $\text{bounded-clinear cblinfun-scale } C\text{-right}$
 $\langle \text{proof} \rangle$

lift-definition $\text{cblinfun-scale } C\text{-left} :: 'a :: \text{complex-normed-vector} \Rightarrow \text{complex} \Rightarrow_{CL} 'a$
is $\lambda x y. y *_C x$
 $\langle \text{proof} \rangle$
lemmas $[simp] = \text{cblinfun-scale } C\text{-left.rep-eq}$

lemma $\text{bounded-clinear-cblinfun-scale } C\text{-left}[\text{bounded-clinear}]$: $\text{bounded-clinear cblinfun-scale } C\text{-left}$
 $\langle \text{proof} \rangle$

lift-definition $\text{cblinfun-mult-right} :: 'a \Rightarrow 'a \Rightarrow_{CL} 'a :: \text{complex-normed-algebra}$ **is** $(*)$
 $\langle \text{proof} \rangle$
declare $\text{cblinfun-mult-right.rep-eq}[simp]$

lemma $\text{bounded-clinear-cblinfun-mult-right}[\text{bounded-clinear}]$: $\text{bounded-clinear cblinfun-mult-right}$
 $\langle \text{proof} \rangle$

lift-definition $\text{cblinfun-mult-left} :: 'a :: \text{complex-normed-algebra} \Rightarrow 'a \Rightarrow_{CL} 'a$ **is** $\lambda x y. y * x$
 $\langle \text{proof} \rangle$
lemmas $[simp] = \text{cblinfun-mult-left.rep-eq}$

lemma $\text{bounded-clinear-cblinfun-mult-left}[\text{bounded-clinear}]$: $\text{bounded-clinear cblinfun-mult-left}$
 $\langle \text{proof} \rangle$

lemmas $\text{bounded-clinear-function-uniform-limit-intros}[\text{uniform-limit-intros}] =$
 $\text{bounded-clinear.uniform-limit[OF bounded-clinear-apply-cblinfun]}$
 $\text{bounded-clinear.uniform-limit[OF bounded-clinear-cblinfun-apply]}$
 $\text{bounded-antilinear.uniform-limit[OF bounded-antilinear-cblinfun-matrix]}$

12.7 The strong operator topology on continuous linear operators

Let $'a$ and $'b$ be two normed real vector spaces. Then the space of linear continuous operators from $'a$ to $'b$ has a canonical norm, and therefore a canonical corresponding topology (the type classes instantiation are given in `Complex_Bounded_Linear_Function0.thy`).

However, there is another topology on this space, the strong operator topology, where T_n tends to T iff, for all x in $'a$, then $T_n x$ tends to $T x$. This is precisely the product topology where the target space is endowed with the norm topology. It is especially useful when $'b$ is the set of real numbers, since then this topology is compact.

We can not implement it using type classes as there is already a topology, but at least we can define it as a topology.

Note that there is yet another (common and useful) topology on operator spaces, the weak operator topology, defined analogously using the product topology, but where the target space is given the weak-* topology, i.e., the pullback of the weak topology on the bidual of the space under the canonical embedding of a space into its bidual. We do not define it there, although it could also be defined analogously.

definition *cstrong-operator-topology*::($'a::\text{complex-normed-vector} \Rightarrow_{CL} 'b::\text{complex-normed-vector}$) topology

where *cstrong-operator-topology* = *pullback-topology UNIV cblinfun-apply euclidean*

lemma *cstrong-operator-topology-topspace*:

topspace cstrong-operator-topology = *UNIV*

<proof>

lemma *cstrong-operator-topology-basis*:

fixes $f::('a::\text{complex-normed-vector} \Rightarrow_{CL} 'b::\text{complex-normed-vector})$ **and** $U::'i \Rightarrow 'b$ set **and** $x::'i \Rightarrow 'a$

assumes *finite I* $\bigwedge i. i \in I \implies \text{open } (U i)$

shows *openin cstrong-operator-topology* $\{f. \forall i \in I. \text{cblinfun-apply } f (x i) \in U i\}$

<proof>

lemma *cstrong-operator-topology-continuous-evaluation*:

continuous-map cstrong-operator-topology euclidean $(\lambda f. \text{cblinfun-apply } f x)$

<proof>

lemma *continuous-on-cstrong-operator-topo-iff-coordinatewise*:

continuous-map T cstrong-operator-topology f

$\longleftrightarrow (\forall x. \text{continuous-map } T \text{ euclidean } (\lambda y. \text{cblinfun-apply } (f y) x))$

<proof>

lemma *cstrong-operator-topology-weaker-than-euclidean*:

```

    continuous-map euclidean cstrong-operator-topology (λf. f)
  ⟨proof⟩
end

```

13 Complex-Bounded-Linear-Function – Complex bounded linear functions (bounded operators)

```

theory Complex-Bounded-Linear-Function

```

```

  imports

```

```

    HOL-Types-To-Sets.Types-To-Sets

```

```

    HOL-Library.Function-Algebras

```

```

    Banach-Steinhaus.Banach-Steinhaus

```

```

    Wlog.Wlog

```

```

    Complex-Inner-Product

```

```

    One-Dimensional-Spaces

```

```

    Complex-Bounded-Linear-Function0

```

```

begin

```

```

unbundle lattice-syntax

```

13.1 Misc basic facts and declarations

```

notation cblinfun-apply (infixr ⟨*V⟩ 70)

```

```

lemma id-cblinfun-apply[simp]: id-cblinfun *V ψ = ψ
  ⟨proof⟩

```

```

lemma apply-id-cblinfun[simp]: ⟨(*V) id-cblinfun = id⟩
  ⟨proof⟩

```

```

lemma isCont-cblinfun-apply[simp]: isCont ((*V) A) ψ
  ⟨proof⟩

```

```

declare cblinfun.scaleC-left[simp]

```

```

lemma cblinfun-apply-clinear[simp]: ⟨clinear (cblinfun-apply A)⟩
  ⟨proof⟩

```

```

lemma cblinfun-cinner-eqI:

```

```

  fixes A B :: ⟨'a::hilbert-space ⇒CL 'a⟩

```

```

  assumes ⟨∧ψ. norm ψ = 1 ⇒ cinner ψ (A *V ψ) = cinner ψ (B *V ψ)⟩

```

```

  shows ⟨A = B⟩

```

```

  ⟨proof⟩

```

```

lemma id-cblinfun-not-0[simp]: ⟨(id-cblinfun :: 'a::{complex-normed-vector, not-singleton}
⇒CL -) ≠ 0⟩
  ⟨proof⟩

```

lemma *cblinfun-norm-geqI*:
assumes $\langle \text{norm } (f *_{\mathcal{V}} x) / \text{norm } x \geq K \rangle$
shows $\langle \text{norm } f \geq K \rangle$
 $\langle \text{proof} \rangle$

declare *scaleC-conv-of-complex[simp]*

lemma *cblinfun-eq-0-on-span*:
fixes $S::\langle 'a::\text{complex-normed-vector set} \rangle$
assumes $x \in \text{cspan } S$
and $\bigwedge s. s \in S \implies F *_{\mathcal{V}} s = 0$
shows $\langle F *_{\mathcal{V}} x = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-eq-on-span*:
fixes $S::\langle 'a::\text{complex-normed-vector set} \rangle$
assumes $x \in \text{cspan } S$
and $\bigwedge s. s \in S \implies F *_{\mathcal{V}} s = G *_{\mathcal{V}} s$
shows $\langle F *_{\mathcal{V}} x = G *_{\mathcal{V}} x \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-eq-0-on-UNIV-span*:
fixes $\text{basis}::\langle 'a::\text{complex-normed-vector set} \rangle$
assumes $\text{cspan basis} = \text{UNIV}$
and $\bigwedge s. s \in \text{basis} \implies F *_{\mathcal{V}} s = 0$
shows $\langle F = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-eq-on-UNIV-span*:
fixes $\text{basis}::'a::\text{complex-normed-vector set}$ **and** $\varphi::'a \Rightarrow 'b::\text{complex-normed-vector}$
assumes $\text{cspan basis} = \text{UNIV}$
and $\bigwedge s. s \in \text{basis} \implies F *_{\mathcal{V}} s = G *_{\mathcal{V}} s$
shows $\langle F = G \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-eq-on-canonical-basis*:
fixes $f g::'a::\{\text{basis-enum}, \text{complex-normed-vector}\} \Rightarrow_{\text{CL}} 'b::\text{complex-normed-vector}$
defines $\text{basis} == \text{set } (\text{canonical-basis}::'a \text{ list})$
assumes $\bigwedge u. u \in \text{basis} \implies f *_{\mathcal{V}} u = g *_{\mathcal{V}} u$
shows $f = g$
 $\langle \text{proof} \rangle$

lemma *cblinfun-eq-0-on-canonical-basis*:
fixes $f::'a::\{\text{basis-enum}, \text{complex-normed-vector}\} \Rightarrow_{\text{CL}} 'b::\text{complex-normed-vector}$
defines $\text{basis} == \text{set } (\text{canonical-basis}::'a \text{ list})$
assumes $\bigwedge u. u \in \text{basis} \implies f *_{\mathcal{V}} u = 0$
shows $f = 0$

⟨proof⟩

lemma *cinner-canonical-basis-eq-0*:

defines $\text{basisA} == \text{set } (\text{canonical-basis}::'a::\text{onb-enum list})$
and $\text{basisB} == \text{set } (\text{canonical-basis}::'b::\text{onb-enum list})$
assumes $\bigwedge u v. u \in \text{basisA} \implies v \in \text{basisB} \implies \text{is-orthogonal } v (F *_V u)$
shows $F = 0$

⟨proof⟩

lemma *cinner-canonical-basis-eq*:

defines $\text{basisA} == \text{set } (\text{canonical-basis}::'a::\text{onb-enum list})$
and $\text{basisB} == \text{set } (\text{canonical-basis}::'b::\text{onb-enum list})$
assumes $\bigwedge u v. u \in \text{basisA} \implies v \in \text{basisB} \implies v \cdot_C (F *_V u) = v \cdot_C (G *_V u)$
shows $F = G$

⟨proof⟩

lemma *cinner-canonical-basis-eq'*:

defines $\text{basisA} == \text{set } (\text{canonical-basis}::'a::\text{onb-enum list})$
and $\text{basisB} == \text{set } (\text{canonical-basis}::'b::\text{onb-enum list})$
assumes $\bigwedge u v. u \in \text{basisA} \implies v \in \text{basisB} \implies (F *_V u) \cdot_C v = (G *_V u) \cdot_C v$
shows $F = G$

⟨proof⟩

lemma *not-not-singleton-cblinfun-zero*:

$\langle x = 0 \rangle$ **if** $\langle \neg \text{class.not-singleton TYPE('a)} \rangle$ **for** $x :: \langle 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'b::\text{complex-normed-vector} \rangle$

⟨proof⟩

lemma *cblinfun-norm-approx-witness*:

fixes $A :: \langle 'a::\{\text{not-singleton, complex-normed-vector}\} \Rightarrow_{CL} 'b::\text{complex-normed-vector} \rangle$
assumes $\langle \varepsilon > 0 \rangle$
shows $\langle \exists \psi. \text{norm } (A *_V \psi) \geq \text{norm } A - \varepsilon \wedge \text{norm } \psi = 1 \rangle$

⟨proof⟩

lemma *cblinfun-norm-approx-witness-mult*:

fixes $A :: \langle 'a::\{\text{not-singleton, complex-normed-vector}\} \Rightarrow_{CL} 'b::\text{complex-normed-vector} \rangle$
assumes $\langle \varepsilon < 1 \rangle$
shows $\langle \exists \psi. \text{norm } (A *_V \psi) \geq \text{norm } A * \varepsilon \wedge \text{norm } \psi = 1 \rangle$

⟨proof⟩

lemma *cblinfun-norm-approx-witness'*:

fixes $A :: \langle 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'b::\text{complex-normed-vector} \rangle$
assumes $\langle \varepsilon > 0 \rangle$
shows $\langle \exists \psi. \text{norm } (A *_V \psi) / \text{norm } \psi \geq \text{norm } A - \varepsilon \rangle$

⟨proof⟩

lemma *cblinfun-to-CARD-1-0[simp]*: $\langle (A :: - \Rightarrow_{CL} -: \text{CARD-1}) = 0 \rangle$

⟨proof⟩

lemma *cblinfun-from-CARD-1-0[simp]*: $\langle (A :: \text{CARD-1} \Rightarrow_{CL} -) = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-cspan-UNIV*:
fixes *basis* :: $\langle ('a :: \{\text{complex-normed-vector}, \text{cfinite-dim}\} \Rightarrow_{CL} 'b :: \text{complex-normed-vector})$
 $\text{set} \rangle$
and *basisA* :: $\langle 'a \text{ set} \rangle$ **and** *basisB* :: $\langle 'b \text{ set} \rangle$
assumes $\langle \text{cspan } \text{basisA} = \text{UNIV} \rangle$ **and** $\langle \text{cspan } \text{basisB} = \text{UNIV} \rangle$
assumes *basis*: $\langle \bigwedge a \ b. a \in \text{basisA} \implies b \in \text{basisB} \implies \exists F \in \text{basis}. \forall a' \in \text{basisA}. F *_{\mathcal{V}}$
 $a' = (\text{if } a' = a \text{ then } b \text{ else } 0) \rangle$
shows $\langle \text{cspan } \text{basis} = \text{UNIV} \rangle$
 $\langle \text{proof} \rangle$

instance *cblinfun* :: $(\langle \{\text{cfinite-dim}, \text{complex-normed-vector}\} \rangle, \langle \{\text{cfinite-dim}, \text{complex-normed-vector}\} \rangle)$
 cfinite-dim
 $\langle \text{proof} \rangle$

lemma *norm-cblinfun-bound-dense*:
assumes $\langle 0 \leq b \rangle$
assumes *S*: $\langle \text{closure } S = \text{UNIV} \rangle$
assumes *bound*: $\langle \bigwedge x. x \in S \implies \text{norm } (\text{cblinfun-apply } f \ x) \leq b * \text{norm } x \rangle$
shows $\langle \text{norm } f \leq b \rangle$
 $\langle \text{proof} \rangle$

lemma *infsum-cblinfun-apply*:
assumes $\langle g \text{ summable-on } S \rangle$
shows $\langle \text{infsum } (\lambda x. A *_{\mathcal{V}} g \ x) \ S = A *_{\mathcal{V}} (\text{infsum } g \ S) \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-cblinfun-apply*:
assumes $\langle (g \text{ has-sum } x) \ S \rangle$
shows $\langle ((\lambda x. A *_{\mathcal{V}} g \ x) \text{ has-sum } (A *_{\mathcal{V}} x)) \ S \rangle$
 $\langle \text{proof} \rangle$

lemma *abs-summable-on-cblinfun-apply*:
assumes $\langle g \text{ abs-summable-on } S \rangle$
shows $\langle (\lambda x. A *_{\mathcal{V}} g \ x) \text{ abs-summable-on } S \rangle$
 $\langle \text{proof} \rangle$

lemma *summable-on-cblinfun-apply*:
assumes $\langle g \text{ summable-on } S \rangle$
shows $\langle (\lambda x. A *_{\mathcal{V}} g \ x) \text{ summable-on } S \rangle$
 $\langle \text{proof} \rangle$

lemma *summable-on-cblinfun-apply-left*:
assumes $\langle A \text{ summable-on } S \rangle$

shows $\langle (\lambda x. A \ x \ *_V \ g) \text{ summable-on } S \rangle$
 $\langle \text{proof} \rangle$

lemma *abs-summable-on-cblinfun-apply-left*:
assumes $\langle A \text{ abs-summable-on } S \rangle$
shows $\langle (\lambda x. A \ x \ *_V \ g) \text{ abs-summable-on } S \rangle$
 $\langle \text{proof} \rangle$

lemma *infsum-cblinfun-apply-left*:
assumes $\langle A \text{ summable-on } S \rangle$
shows $\langle \text{infsum } (\lambda x. A \ x \ *_V \ g) \ S = (\text{infsum } A \ S) \ *_V \ g \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-cblinfun-apply-left*:
assumes $\langle (A \text{ has-sum } x) \ S \rangle$
shows $\langle ((\lambda x. A \ x \ *_V \ g) \text{ has-sum } (x \ *_V \ g)) \ S \rangle$
 $\langle \text{proof} \rangle$

The next eight lemmas logically belong in *Complex-Bounded-Operators.Complex-Inner-Product* but the proofs use facts from this theory.

lemma *has-sum-cinner-left*:
assumes $\langle (f \text{ has-sum } x) \ I \rangle$
shows $\langle ((\lambda i. \text{cinner } a \ (f \ i)) \text{ has-sum } \text{cinner } a \ x) \ I \rangle$
 $\langle \text{proof} \rangle$

lemma *summable-on-cinner-left*:
assumes $\langle f \text{ summable-on } I \rangle$
shows $\langle (\lambda i. \text{cinner } a \ (f \ i)) \text{ summable-on } I \rangle$
 $\langle \text{proof} \rangle$

lemma *infsum-cinner-left*:
assumes $\langle \varphi \text{ summable-on } I \rangle$
shows $\langle \text{cinner } \psi \ (\sum_{\infty i \in I. \varphi \ i}) = (\sum_{\infty i \in I. \text{cinner } \psi \ (\varphi \ i)) \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-cinner-right*:
assumes $\langle (f \text{ has-sum } x) \ I \rangle$
shows $\langle ((\lambda i. f \ i \ \cdot_C \ a) \text{ has-sum } (x \ \cdot_C \ a)) \ I \rangle$
 $\langle \text{proof} \rangle$

lemma *summable-on-cinner-right*:
assumes $\langle f \text{ summable-on } I \rangle$
shows $\langle (\lambda i. f \ i \ \cdot_C \ a) \text{ summable-on } I \rangle$
 $\langle \text{proof} \rangle$

lemma *infsum-cinner-right*:
assumes $\langle \varphi \text{ summable-on } I \rangle$
shows $\langle (\sum_{\infty i \in I. \varphi \ i} \cdot_C \ \psi = (\sum_{\infty i \in I. \varphi \ i \ \cdot_C \ \psi}) \rangle$
 $\langle \text{proof} \rangle$

lemma *Cauchy-cinner-product-summable*:
assumes *asum*: $\langle a \text{ summable-on } UNIV \rangle$
assumes *bsum*: $\langle b \text{ summable-on } UNIV \rangle$
assumes $\langle \text{finite } X \rangle \langle \text{finite } Y \rangle$
assumes *pos*: $\langle \bigwedge x y. x \notin X \implies y \notin Y \implies \text{cinner } (a \ x) \ (b \ y) \geq 0 \rangle$
shows *absum*: $\langle (\lambda(x, y). \text{cinner } (a \ x) \ (b \ y)) \text{ summable-on } UNIV \rangle$
 $\langle \text{proof} \rangle$

A variant of *Series.Cauchy-product-sums* with $(*)$ replaced by $(\cdot_C)$. Differently from *Series.Cauchy-product-sums*, we do not require absolute summability of a and b individually but only unconditional summability of a , b , and their product. While on, e.g., reals, unconditional summability is equivalent to absolute summability, in general unconditional summability is a weaker requirement.

Logically belong in *Complex-Bounded-Operators.Complex-Inner-Product* but the proof uses facts from this theory.

lemma
fixes $a \ b :: \text{nat} \Rightarrow 'a :: \text{complex-inner}$
assumes *asum*: $\langle a \text{ summable-on } UNIV \rangle$
assumes *bsum*: $\langle b \text{ summable-on } UNIV \rangle$
assumes *absum*: $\langle (\lambda(x, y). \text{cinner } (a \ x) \ (b \ y)) \text{ summable-on } UNIV \rangle$

shows *Cauchy-cinner-product-infsum*: $\langle (\sum_{\infty} k. \sum_{i \leq k}. \text{cinner } (a \ i) \ (b \ (k - i))) = \text{cinner } (\sum_{\infty} k. a \ k) \ (\sum_{\infty} k. b \ k) \rangle$
and *Cauchy-cinner-product-infsum-exists*: $\langle (\lambda k. \sum_{i \leq k}. \text{cinner } (a \ i) \ (b \ (k - i))) \text{ summable-on } UNIV \rangle$
 $\langle \text{proof} \rangle$

lemma *CBlinfun-plus*:
assumes [*simp*]: $\langle \text{bounded-clinear } f \rangle \langle \text{bounded-clinear } g \rangle$
shows $\langle \text{CBlinfun } (f + g) = \text{CBlinfun } f + \text{CBlinfun } g \rangle$
 $\langle \text{proof} \rangle$

lemma *CBlinfun-scaleC*:
assumes $\langle \text{bounded-clinear } f \rangle$
shows $\langle \text{CBlinfun } (\lambda y. c *_{\mathbb{C}} f \ y) = c *_{\mathbb{C}} \text{CBlinfun } f \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-sup-norm-cblinfun*:
fixes $A :: \langle 'a :: \{\text{complex-normed-vector}, \text{not-singleton}\} \Rightarrow_{\mathbb{C}L} 'b :: \text{complex-inner} \rangle$
shows $\langle \text{norm } A = (\text{SUP } (\psi, \varphi). \text{cmod } (\text{cinner } \psi \ (A *_{\mathbb{V}} \varphi)) / (\text{norm } \psi * \text{norm } \varphi)) \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-cblinfun-Sup*: $\langle \text{norm } A = (\text{SUP } \psi. \text{norm } (A *_{\mathbb{V}} \psi) / \text{norm } \psi) \rangle$

⟨proof⟩

lemma *cblinfun-eq-on*:

fixes $A\ B :: 'a::\text{cbanach} \Rightarrow_{CL} 'b::\text{complex-normed-vector}$
assumes $\bigwedge x. x \in G \implies A *_V x = B *_V x$ **and** $\langle t \in \text{closure } (\text{cspan } G) \rangle$
shows $A *_V t = B *_V t$
 ⟨proof⟩

lemma *cblinfun-eq-gen-eqI*:

fixes $A\ B :: 'a::\text{cbanach} \Rightarrow_{CL} 'b::\text{complex-normed-vector}$
assumes $\bigwedge x. x \in G \implies A *_V x = B *_V x$ **and** $\langle \text{ccspan } G = \top \rangle$
shows $A = B$
 ⟨proof⟩

declare *cnj-bounded-antilinear*[*bounded-antilinear*]

lemma *Cblinfun-comp-bounded-cbilinear*: $\langle \text{bounded-clinear } (CBlinfun\ o\ p) \rangle$ **if** $\langle \text{bounded-cbilinear } p \rangle$
 ⟨proof⟩

lemma *Cblinfun-comp-bounded-sesquilinear*: $\langle \text{bounded-antilinear } (CBlinfun\ o\ p) \rangle$
if $\langle \text{bounded-sesquilinear } p \rangle$
 ⟨proof⟩

declare (**in** *bounded-cbilinear*) *bounded-cbilinear-axioms*[*bounded-cbilinear*]
declare (**in** *bounded-clinear*) *bounded-clinear-axioms*[*bounded-clinear*]

lemma *norm-cblinfun-bound-both-sides*:

fixes $a :: \langle 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'b::\text{complex-inner} \rangle$
assumes $\langle b \geq 0 \rangle$
assumes *leg*: $\langle \bigwedge \psi\ \varphi. \text{norm } \psi = 1 \implies \text{norm } \varphi = 1 \implies \text{norm } (\psi \cdot_C a\ \varphi) \leq b \rangle$
shows $\langle \text{norm } a \leq b \rangle$
 ⟨proof⟩

lemma *norm-cblinfun-bound-unit*:

assumes $\langle b \geq 0 \rangle$
assumes $\langle \bigwedge \psi. \text{norm } \psi = 1 \implies \text{norm } (a *_V \psi) \leq b \rangle$
shows $\langle \text{norm } a \leq b \rangle$
 ⟨proof⟩

lemma *cblinfun-eq-from-separatingI*:

fixes $a\ b :: \langle 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'b::\text{complex-normed-vector} \rangle$
assumes $\langle \text{separating-set } (\text{bounded-clinear} :: ('a \Rightarrow 'b) \Rightarrow \text{bool})\ S \rangle$
assumes $\langle \bigwedge x. x \in S \implies a\ x = b\ x \rangle$
shows $\langle a = b \rangle$
 ⟨proof⟩

lemma *cblinfun-eq-from-separatingI2*:

fixes $a\ b :: \langle 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'b::\text{complex-normed-vector} \rangle$

```

assumes ‹separating-set (bounded-clinear :: ('a ⇒ 'b) ⇒ bool) ((λ(x,y). h x y) ‹
(S × T))›
assumes ‹∧x y. x ∈ S ⇒ y ∈ T ⇒ a (h x y) = b (h x y)›
shows ‹a = b›
‹proof›

```

13.2 Relationship to real bounded operators ($- \Rightarrow_L -$)

instantiation *blinfun* :: (*real-normed-vector*, *complex-normed-vector*) *complex-normed-vector*
begin

lift-definition *scaleC-blinfun* :: ‹*complex* ⇒
('a::*real-normed-vector*, 'b::*complex-normed-vector*) *blinfun* ⇒
('a, 'b) *blinfun*›
is ‹λ *c*::*complex*. λ *f*::'a⇒'b. (λ *x*. *c* *_{*C*} (*f* *x*))›
‹*proof*›

instance
‹*proof*›
end

lemma *clinear-blinfun-compose-left*: ‹*clinear* (λ*x*. *blinfun-compose* *x* *y*)›
‹*proof*›

instance *blinfun* :: (*real-normed-vector*, *cbanach*) *cbanach*‹*proof*›

lemma *blinfun-compose-assoc*: (*A* *o_L* *B*) *o_L* *C* = *A* *o_L* (*B* *o_L* *C*)
‹*proof*›

lift-definition *blinfun-of-cblinfun*::‹'a::*complex-normed-vector* ⇒_{*C_L*} 'b::*complex-normed-vector*
⇒ 'a ⇒_{*L*} 'b› **is** *id*
‹*proof*›

lift-definition *blinfun-cblinfun-eq* ::
‹'a ⇒_{*L*} 'b ⇒ 'a::*complex-normed-vector* ⇒_{*C_L*} 'b::*complex-normed-vector* ⇒ bool›
is (=) ‹*proof*›

lemma *blinfun-cblinfun-eq-bi-unique*[*transfer-rule*]: ‹*bi-unique* *blinfun-cblinfun-eq*›
‹*proof*›

lemma *blinfun-cblinfun-eq-right-total*[*transfer-rule*]: ‹*right-total* *blinfun-cblinfun-eq*›
‹*proof*›

named-theorems *cblinfun-blinfun-transfer*

lemma *cblinfun-blinfun-transfer-0*[*cblinfun-blinfun-transfer*]:
blinfun-cblinfun-eq (0::(-,-) *blinfun*) (0::(-,-) *cblinfun*)
‹*proof*›

lemma *cblinfun-blinfun-transfer-plus*[*cblinfun-blinfun-transfer*]:
includes *lifting-syntax*
shows (*blinfun-cblinfun-eq* $\implies$ *blinfun-cblinfun-eq* $\implies$ *blinfun-cblinfun-eq*)
 (+) (+)
 ⟨*proof*⟩

lemma *cblinfun-blinfun-transfer-minus*[*cblinfun-blinfun-transfer*]:
includes *lifting-syntax*
shows (*blinfun-cblinfun-eq* $\implies$ *blinfun-cblinfun-eq* $\implies$ *blinfun-cblinfun-eq*)
 (−) (−)
 ⟨*proof*⟩

lemma *cblinfun-blinfun-transfer-uminus*[*cblinfun-blinfun-transfer*]:
includes *lifting-syntax*
shows (*blinfun-cblinfun-eq* $\implies$ *blinfun-cblinfun-eq*) (*uminus*) (*uminus*)
 ⟨*proof*⟩

definition *real-complex-eq* *r c* $\longleftrightarrow$ *complex-of-real* *r* = *c*

lemma *bi-unique-real-complex-eq*[*transfer-rule*]: ⟨*bi-unique real-complex-eq*⟩
 ⟨*proof*⟩

lemma *left-total-real-complex-eq*[*transfer-rule*]: ⟨*left-total real-complex-eq*⟩
 ⟨*proof*⟩

lemma *cblinfun-blinfun-transfer-scaleC*[*cblinfun-blinfun-transfer*]:
includes *lifting-syntax*
shows (*real-complex-eq* $\implies$ *blinfun-cblinfun-eq* $\implies$ *blinfun-cblinfun-eq*)
 (*scaleR*) (*scaleC*)
 ⟨*proof*⟩

lemma *cblinfun-blinfun-transfer-CBlinfun*[*cblinfun-blinfun-transfer*]:
includes *lifting-syntax*
shows (*eq-onp bounded-clinear* $\implies$ *blinfun-cblinfun-eq*) *Blinfun* *CBlinfun*
 ⟨*proof*⟩

lemma *cblinfun-blinfun-transfer-norm*[*cblinfun-blinfun-transfer*]:
includes *lifting-syntax*
shows (*blinfun-cblinfun-eq* $\implies$ (=)) *norm* *norm*
 ⟨*proof*⟩

lemma *cblinfun-blinfun-transfer-dist*[*cblinfun-blinfun-transfer*]:
includes *lifting-syntax*
shows (*blinfun-cblinfun-eq* $\implies$ *blinfun-cblinfun-eq* $\implies$ (=)) *dist* *dist*
 ⟨*proof*⟩

lemma *cblinfun-blinfun-transfer-sgn*[*cblinfun-blinfun-transfer*]:
includes *lifting-syntax*
shows (*blinfun-cblinfun-eq* $\implies$ *blinfun-cblinfun-eq*) *sgn* *sgn*

$\langle \text{proof} \rangle$

lemma *cblinfun-blinfun-transfer-Cauchy*[*cblinfun-blinfun-transfer*]:
 includes *lifting-syntax*
 shows $((=) \implies \text{blinfun-cblinfun-eq}) \implies (=)$ *Cauchy Cauchy*
 $\langle \text{proof} \rangle$

lemma *cblinfun-blinfun-transfer-tendsto*[*cblinfun-blinfun-transfer*]:
 includes *lifting-syntax*
 shows $((=) \implies \text{blinfun-cblinfun-eq}) \implies \text{blinfun-cblinfun-eq} \implies (=)$
 $\implies (=)$ *tendsto tendsto*
 $\langle \text{proof} \rangle$

lemma *cblinfun-blinfun-transfer-compose*[*cblinfun-blinfun-transfer*]:
 includes *lifting-syntax*
 shows $(\text{blinfun-cblinfun-eq} \implies \text{blinfun-cblinfun-eq} \implies \text{blinfun-cblinfun-eq})$
 $(o_L) (o_{CL})$
 $\langle \text{proof} \rangle$

lemma *cblinfun-blinfun-transfer-apply*[*cblinfun-blinfun-transfer*]:
 includes *lifting-syntax*
 shows $(\text{blinfun-cblinfun-eq} \implies (=) \implies (=))$ *blinfun-apply cblinfun-apply*
 $\langle \text{proof} \rangle$

lemma *blinfun-of-cblinfun-inj*:
 $\langle \text{blinfun-of-cblinfun } f = \text{blinfun-of-cblinfun } g \implies f = g \rangle$
 $\langle \text{proof} \rangle$

lemma *blinfun-of-cblinfun-inv*:
 assumes $\bigwedge c. \bigwedge x. f *_v (c *_C x) = c *_C (f *_v x)$
 shows $\exists g. \text{blinfun-of-cblinfun } g = f$
 $\langle \text{proof} \rangle$

lemma *blinfun-of-cblinfun-zero*:
 $\langle \text{blinfun-of-cblinfun } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *blinfun-of-cblinfun-uminus*:
 $\langle \text{blinfun-of-cblinfun } (- f) = - (\text{blinfun-of-cblinfun } f) \rangle$
 $\langle \text{proof} \rangle$

lemma *blinfun-of-cblinfun-minus*:
 $\langle \text{blinfun-of-cblinfun } (f - g) = \text{blinfun-of-cblinfun } f - \text{blinfun-of-cblinfun } g \rangle$
 $\langle \text{proof} \rangle$

lemma *blinfun-of-cblinfun-scaleC*:
 $\langle \text{blinfun-of-cblinfun } (c *_C f) = c *_C (\text{blinfun-of-cblinfun } f) \rangle$
 $\langle \text{proof} \rangle$

lemma *blinfun-of-cblinfun-scaleR*:

$\langle \text{blinfun-of-cblinfun } (c *_{\mathbb{R}} f) = c *_{\mathbb{R}} (\text{blinfun-of-cblinfun } f) \rangle$

$\langle \text{proof} \rangle$

lemma *blinfun-of-cblinfun-norm*:

fixes $f :: \langle 'a :: \text{complex-normed-vector} \Rightarrow_{CL} 'b :: \text{complex-normed-vector} \rangle$

shows $\langle \text{norm } f = \text{norm } (\text{blinfun-of-cblinfun } f) \rangle$

$\langle \text{proof} \rangle$

lemma *blinfun-of-cblinfun-cblinfun-compose*:

fixes $f :: \langle 'b :: \text{complex-normed-vector} \Rightarrow_{CL} 'c :: \text{complex-normed-vector} \rangle$

and $g :: \langle 'a :: \text{complex-normed-vector} \Rightarrow_{CL} 'b \rangle$

shows $\langle \text{blinfun-of-cblinfun } (f \circ_{CL} g) = (\text{blinfun-of-cblinfun } f) \circ_L (\text{blinfun-of-cblinfun } g) \rangle$

$\langle \text{proof} \rangle$

13.3 Composition

lemma *cblinfun-compose-assoc*:

shows $\langle (A \circ_{CL} B) \circ_{CL} C = A \circ_{CL} (B \circ_{CL} C) \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-compose-zero-right[simp]*: $U \circ_{CL} 0 = 0$

$\langle \text{proof} \rangle$

lemma *cblinfun-compose-zero-left[simp]*: $0 \circ_{CL} U = 0$

$\langle \text{proof} \rangle$

lemma *cblinfun-compose-scaleC-left[simp]*:

fixes $A :: \langle 'b :: \text{complex-normed-vector} \Rightarrow_{CL} 'c :: \text{complex-normed-vector} \rangle$

and $B :: \langle 'a :: \text{complex-normed-vector} \Rightarrow_{CL} 'b \rangle$

shows $\langle (a *_{\mathbb{C}} A) \circ_{CL} B = a *_{\mathbb{C}} (A \circ_{CL} B) \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-compose-scaleR-left[simp]*:

fixes $A :: \langle 'b :: \text{complex-normed-vector} \Rightarrow_{CL} 'c :: \text{complex-normed-vector} \rangle$

and $B :: \langle 'a :: \text{complex-normed-vector} \Rightarrow_{CL} 'b \rangle$

shows $\langle (a *_{\mathbb{R}} A) \circ_{CL} B = a *_{\mathbb{R}} (A \circ_{CL} B) \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-compose-scaleC-right[simp]*:

fixes $A :: \langle 'b :: \text{complex-normed-vector} \Rightarrow_{CL} 'c :: \text{complex-normed-vector} \rangle$

and $B :: \langle 'a :: \text{complex-normed-vector} \Rightarrow_{CL} 'b \rangle$

shows $\langle A \circ_{CL} (a *_{\mathbb{C}} B) = a *_{\mathbb{C}} (A \circ_{CL} B) \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-compose-scaleR-right[simp]*:

fixes $A :: \langle 'b :: \text{complex-normed-vector} \Rightarrow_{CL} 'c :: \text{complex-normed-vector} \rangle$

and $B :: \langle 'a :: \text{complex-normed-vector} \Rightarrow_{CL} 'b \rangle$

shows $\langle A \circ_{CL} (a *_R B) = a *_R (A \circ_{CL} B) \rangle$
 $\langle proof \rangle$

lemma *cblinfun-compose-id-right*[simp]:
shows $U \circ_{CL} id\text{-}cblinfun = U$
 $\langle proof \rangle$

lemma *cblinfun-compose-id-left*[simp]:
shows $id\text{-}cblinfun \circ_{CL} U = U$
 $\langle proof \rangle$

lemma *cblinfun-compose-add-left*: $\langle (a + b) \circ_{CL} c = (a \circ_{CL} c) + (b \circ_{CL} c) \rangle$
 $\langle proof \rangle$

lemma *cblinfun-compose-add-right*: $\langle a \circ_{CL} (b + c) = (a \circ_{CL} b) + (a \circ_{CL} c) \rangle$
 $\langle proof \rangle$

lemma *cbilinear-cblinfun-compose*[simp]: *cbilinear cblinfun-compose*
 $\langle proof \rangle$

lemma *cblinfun-compose-sum-left*: $\langle (\sum_{i \in S} g \ i) \circ_{CL} x = (\sum_{i \in S} g \ i \circ_{CL} x) \rangle$
 $\langle proof \rangle$

lemma *cblinfun-compose-sum-right*: $\langle x \circ_{CL} (\sum_{i \in S} g \ i) = (\sum_{i \in S} x \circ_{CL} g \ i) \rangle$
 $\langle proof \rangle$

lemma *cblinfun-compose-minus-right*: $\langle a \circ_{CL} (b - c) = (a \circ_{CL} b) - (a \circ_{CL} c) \rangle$
 $\langle proof \rangle$

lemma *cblinfun-compose-minus-left*: $\langle (a - b) \circ_{CL} c = (a \circ_{CL} c) - (b \circ_{CL} c) \rangle$
 $\langle proof \rangle$

lemma *simp-a-oCL-b*: $\langle a \circ_{CL} b = c \implies a \circ_{CL} (b \circ_{CL} d) = c \circ_{CL} d \rangle$
— A convenience lemma to transform simplification rules of the form $a \circ_{CL} b = c$. E.g., *simp-a-oCL-b*[*OF isometryD, simp*] declares a simp-rule for simplifying $adj \ x \circ_{CL} (x \circ_{CL} y) = id\text{-}cblinfun \circ_{CL} y$.
 $\langle proof \rangle$

lemma *simp-a-oCL-b'*: $\langle a \circ_{CL} b = c \implies a *_V (b *_V d) = c *_V d \rangle$
— A convenience lemma to transform simplification rules of the form $a \circ_{CL} b = c$. E.g., *simp-a-oCL-b'*[*OF isometryD, simp*] declares a simp-rule for simplifying $adj \ x *_V x *_V y = id\text{-}cblinfun *_V y$.
 $\langle proof \rangle$

lemma *cblinfun-compose-uminus-left*: $\langle (- \ a) \circ_{CL} b = - \ (a \circ_{CL} b) \rangle$
 $\langle proof \rangle$

lemma *cblinfun-compose-uminus-right*: $\langle a \circ_{CL} (- \ b) = - \ (a \circ_{CL} b) \rangle$
 $\langle proof \rangle$

lemma *bounded-clinear-cblinfun-compose-left*: $\langle \text{bounded-clinear } (\lambda x. x \text{ } o_{CL} \text{ } y) \rangle$
 $\langle \text{proof} \rangle$
lemma *bounded-clinear-cblinfun-compose-right*: $\langle \text{bounded-clinear } (\lambda y. x \text{ } o_{CL} \text{ } y) \rangle$
 $\langle \text{proof} \rangle$
lemma *clinear-cblinfun-compose-left*: $\langle \text{clinear } (\lambda x. x \text{ } o_{CL} \text{ } y) \rangle$
 $\langle \text{proof} \rangle$
lemma *clinear-cblinfun-compose-right*: $\langle \text{clinear } (\lambda y. x \text{ } o_{CL} \text{ } y) \rangle$
 $\langle \text{proof} \rangle$

lemma *additive-cblinfun-compose-left[simp]*: $\langle \text{Modules.additive } (\lambda x. x \text{ } o_{CL} \text{ } a) \rangle$
 $\langle \text{proof} \rangle$
lemma *additive-cblinfun-compose-right[simp]*: $\langle \text{Modules.additive } (\lambda x. a \text{ } o_{CL} \text{ } x) \rangle$
 $\langle \text{proof} \rangle$
lemma *isCont-cblinfun-compose-left*: $\langle \text{isCont } (\lambda x. x \text{ } o_{CL} \text{ } a) \text{ } y \rangle$
 $\langle \text{proof} \rangle$
lemma *isCont-cblinfun-compose-right*: $\langle \text{isCont } (\lambda x. a \text{ } o_{CL} \text{ } x) \text{ } y \rangle$
 $\langle \text{proof} \rangle$

lemma *cspan-compose-closed*:
assumes $\langle \bigwedge a \ b. a \in A \implies b \in A \implies a \text{ } o_{CL} \text{ } b \in A \rangle$
assumes $\langle a \in \text{cspan } A \rangle$ **and** $\langle b \in \text{cspan } A \rangle$
shows $\langle a \text{ } o_{CL} \text{ } b \in \text{cspan } A \rangle$
 $\langle \text{proof} \rangle$

13.4 Adjoint

lift-definition

$\text{adj} :: 'a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \text{complex-inner} \Rightarrow 'b \Rightarrow_{CL} 'a \text{ } (\langle - \rangle [99] \text{ } 100)$
is *cadjoint* $\langle \text{proof} \rangle$

definition *selfadjoint* :: $\langle ('a :: \text{chilbert-space} \Rightarrow_{CL} 'a) \Rightarrow \text{bool} \rangle$ **where**
 $\langle \text{selfadjoint } a \iff a^* = a \rangle$

lemma *id-cblinfun-adjoint[simp]*: $\text{id-cblinfun}^* = \text{id-cblinfun}$
 $\langle \text{proof} \rangle$

lemma *double-adj[simp]*: $(A^*)^* = A$
 $\langle \text{proof} \rangle$

lemma *adj-cblinfun-compose[simp]*:
fixes $B :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \text{chilbert-space} \rangle$
and $A :: \langle 'b \Rightarrow_{CL} 'c :: \text{complex-inner} \rangle$
shows $(A \text{ } o_{CL} \text{ } B)^* = (B^*) \text{ } o_{CL} \text{ } (A^*)$
 $\langle \text{proof} \rangle$

lemma *scaleC-adj[simp]*: $(a *_C A)^* = (\text{cnj } a) *_C (A^*)$
 $\langle \text{proof} \rangle$

lemma *scaleR-adj[simp]*: $(a *_R A)* = a *_R (A*)$
 $\langle \text{proof} \rangle$

lemma *adj-plus*: $\langle (A + B)* = (A*) + (B*) \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-adj-left*:
fixes $G :: 'b::\text{hilbert-space} \Rightarrow_{CL} 'a::\text{complex-inner}$
shows $\langle (G* *_V x) \cdot_C y = x \cdot_C (G *_V y) \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-adj-right*:
fixes $G :: 'b::\text{hilbert-space} \Rightarrow_{CL} 'a::\text{complex-inner}$
shows $\langle x \cdot_C (G* *_V y) = (G *_V x) \cdot_C y \rangle$
 $\langle \text{proof} \rangle$

lemma *adj-0[simp]*: $\langle 0* = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *selfadjoint-0[simp]*: $\langle \text{selfadjoint } 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-adj[simp]*: $\langle \text{norm } (A*) = \text{norm } A \rangle$
for $A :: \langle 'b::\text{hilbert-space} \Rightarrow_{CL} 'c::\text{complex-inner} \rangle$
 $\langle \text{proof} \rangle$

lemma *antilinear-adj[simp]*: $\langle \text{antilinear } \text{adj} \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-antilinear-adj[bounded-antilinear, simp]*: $\langle \text{bounded-antilinear } \text{adj} \rangle$
 $\langle \text{proof} \rangle$

lemma *adjoint-eqI*:
fixes $G :: \langle 'b::\text{hilbert-space} \Rightarrow_{CL} 'a::\text{complex-inner} \rangle$
and $F :: \langle 'a \Rightarrow_{CL} 'b \rangle$
assumes $\langle \bigwedge x y. ((\text{cbfun-apply } F) x \cdot_C y) = (x \cdot_C (\text{cbfun-apply } G) y) \rangle$
shows $\langle F = G* \rangle$
 $\langle \text{proof} \rangle$

lemma *adj-uminus*: $\langle (-A)* = - (A*) \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-real-selfadjointI*:
— Prop. II.2.12 in [1]
assumes $\langle \bigwedge \psi. \psi \cdot_C (A *_V \psi) \in \mathbb{R} \rangle$
shows $\langle \text{selfadjoint } A \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-AAadj[simp]*: $\langle \text{norm } (A \circ_{CL} A^*) = (\text{norm } A)^2 \rangle$ **for** $A :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \{\text{complex-inner}\} \rangle$
 $\langle \text{proof} \rangle$

lemma *sum-adj*: $\langle (\text{sum } a \ F)^* = \text{sum } (\lambda i. (a \ i)^*) \ F \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-adj*:
assumes $\langle f \ \text{has-sum } x \rangle \ I$
shows $\langle ((\lambda x. \ \text{adj } (f \ x)) \ \text{has-sum adj } x) \ I \rangle$
 $\langle \text{proof} \rangle$

lemma *adj-minus*: $\langle (A - B)^* = (A^*) - (B^*) \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-selfadjoint-real*: $\langle x \cdot_C (A *_V x) \in \mathbb{R} \rangle$ **if** $\langle \text{selfadjoint } A \rangle$
 $\langle \text{proof} \rangle$

lemma *adj-inject*: $\langle \text{adj } a = \text{adj } b \longleftrightarrow a = b \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-AadjA[simp]*: $\langle \text{norm } (A^* \circ_{CL} A) = (\text{norm } A)^2 \rangle$ **for** $A :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \text{chilbert-space} \rangle$
 $\langle \text{proof} \rangle$

lemma *cspan-adj-closed*:
assumes $\langle \bigwedge a. a \in A \implies a^* \in A \rangle$
assumes $\langle a \in \text{cspan } A \rangle$
shows $\langle a^* \in \text{cspan } A \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-norm-is-Sup-cinner*:
— [1], Proposition II.2.13
fixes $A :: \langle 'a :: \{\text{not-singleton}, \text{chilbert-space}\} \Rightarrow_{CL} 'a \rangle$
assumes $A \text{selfadj}$: $\langle \text{selfadjoint } A \rangle$
shows $\langle \text{is-Sup } ((\lambda \psi. \text{cmod } (\psi \cdot_C (A *_V \psi))) \ ' \ \{\psi. \text{norm } \psi = 1\}) \ (\text{norm } A) \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-norm-approx-witness-cinner*:
fixes $A :: \langle 'a :: \{\text{not-singleton}, \text{chilbert-space}\} \Rightarrow_{CL} 'a \rangle$
assumes $\langle \text{selfadjoint } A \rangle$ **and** $\langle \varepsilon > 0 \rangle$
shows $\langle \exists \psi. \text{cmod } (\psi \cdot_C (A *_V \psi)) \geq \text{norm } A - \varepsilon \wedge \text{norm } \psi = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-norm-approx-witness-cinner'*:
fixes $A :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'a \rangle$
assumes $\langle \text{selfadjoint } A \rangle$ **and** $\langle \varepsilon > 0 \rangle$

shows $\langle \exists \psi. \text{cmod } (\psi \cdot_C A \psi) / (\text{norm } \psi)^2 \geq \text{norm } A - \varepsilon \rangle$
 $\langle \text{proof} \rangle$

13.5 Powers of operators

lift-definition $\text{cblinfun-power} :: \langle 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'a \Rightarrow \text{nat} \Rightarrow 'a \Rightarrow_{CL} 'a \rangle$ **is**
 $\langle \lambda(a::'a \Rightarrow 'a) n. a \hat{\sim} n \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{cblinfun-power-0}[\text{simp}]$: $\langle \text{cblinfun-power } A \ 0 = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{cblinfun-power-Suc'}$: $\langle \text{cblinfun-power } A \ (\text{Suc } n) = A \ o_{CL} \ \text{cblinfun-power } A \ n \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{cblinfun-power-Suc}$: $\langle \text{cblinfun-power } A \ (\text{Suc } n) = \text{cblinfun-power } A \ n \ o_{CL} \ A \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{cblinfun-power-compose}[\text{simp}]$: $\langle \text{cblinfun-power } A \ n \ o_{CL} \ \text{cblinfun-power } A \ m = \text{cblinfun-power } A \ (n+m) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{cblinfun-power-scaleC}$: $\langle \text{cblinfun-power } (c *_C a) \ n = c \hat{\sim} n *_C \text{cblinfun-power } a \ n \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{cblinfun-power-scaleR}$: $\langle \text{cblinfun-power } (c *_R a) \ n = c \hat{\sim} n *_R \text{cblinfun-power } a \ n \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{cblinfun-power-uminus}$: $\langle \text{cblinfun-power } (-a) \ n = (-1) \hat{\sim} n *_R \text{cblinfun-power } a \ n \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{cblinfun-power-adj}$: $\langle (\text{cblinfun-power } S \ n)^* = \text{cblinfun-power } (S^*) \ n \rangle$
 $\langle \text{proof} \rangle$

13.6 Unitaries / isometries

definition $\text{isometry} :: \langle 'a::\text{hilbert-space} \Rightarrow_{CL} 'b::\text{complex-inner} \Rightarrow \text{bool} \rangle$ **where**
 $\langle \text{isometry } U \longleftrightarrow U^* \ o_{CL} \ U = \text{id-cblinfun} \rangle$

definition $\text{unitary} :: \langle 'a::\text{hilbert-space} \Rightarrow_{CL} 'b::\text{complex-inner} \Rightarrow \text{bool} \rangle$ **where**
 $\langle \text{unitary } U \longleftrightarrow (U^* \ o_{CL} \ U = \text{id-cblinfun}) \wedge (U \ o_{CL} \ U^* = \text{id-cblinfun}) \rangle$

lemma unitaryI : $\langle \text{unitary } a \rangle$ **if** $\langle a^* \ o_{CL} \ a = \text{id-cblinfun} \rangle$ **and** $\langle a \ o_{CL} \ a^* = \text{id-cblinfun} \rangle$

$\langle \text{proof} \rangle$

lemma *unitary-twosided-isometry*: $\text{unitary } U \longleftrightarrow \text{isometry } U \wedge \text{isometry } (U^*)$
 $\langle \text{proof} \rangle$

lemma *isometryD[simp]*: $\text{isometry } U \implies U^* \circ_{CL} U = \text{id-cblinfun}$
 $\langle \text{proof} \rangle$

lemma *unitaryD1*: $\text{unitary } U \implies U^* \circ_{CL} U = \text{id-cblinfun}$
 $\langle \text{proof} \rangle$

lemma *unitaryD2[simp]*: $\text{unitary } U \implies U \circ_{CL} U^* = \text{id-cblinfun}$
 $\langle \text{proof} \rangle$

lemma *unitary-isometry[simp]*: $\text{unitary } U \implies \text{isometry } U$
 $\langle \text{proof} \rangle$

lemma *unitary-adj[simp]*: $\text{unitary } (U^*) = \text{unitary } U$
 $\langle \text{proof} \rangle$

lemma *isometry-cblinfun-compose[simp]*:
 assumes *isometry A and isometry B*
 shows *isometry (A $\circ_{CL}$ B)*
 $\langle \text{proof} \rangle$

lemma *unitary-cblinfun-compose[simp]*: *unitary (A $\circ_{CL}$ B)*
 if *unitary A and unitary B*
 $\langle \text{proof} \rangle$

lemma *unitary-surj*:
 assumes *unitary U*
 shows *surj (cblinfun-apply U)*
 $\langle \text{proof} \rangle$

lemma *unitary-id[simp]*: *unitary id-cblinfun*
 $\langle \text{proof} \rangle$

lemma *orthogonal-on-basis-is-isometry*:
 assumes *spanB: $\langle \text{ccspan } B = \top \rangle$*
 assumes *orthoU: $\langle \bigwedge b \in B. b \in B \implies c \in B \implies \text{cinner } (U *_V b) (U *_V c) = \text{cinner } b \ c \rangle$*
 shows *$\langle \text{isometry } U \rangle$*
 $\langle \text{proof} \rangle$

lemma *isometry-preserves-norm*: *$\langle \text{isometry } U \implies \text{norm } (U *_V \psi) = \text{norm } \psi \rangle$*
 $\langle \text{proof} \rangle$

lemma *norm-isometry-compose*:

assumes $\langle isometry\ U \rangle$
shows $\langle norm\ (U\ o_{CL}\ A) = norm\ A \rangle$
 $\langle proof \rangle$

lemma *norm-isometry*:
fixes $U :: \langle 'a :: \{chilbert-space, not-singleton\} \Rightarrow_{CL}\ 'b :: complex-inner \rangle$
assumes $\langle isometry\ U \rangle$
shows $\langle norm\ U = 1 \rangle$
 $\langle proof \rangle$

lemma *norm-preserving-isometry*: $\langle isometry\ U \rangle$ **if** $\langle \bigwedge \psi. norm\ (U\ *_V\ \psi) = norm\ \psi \rangle$
 $\langle proof \rangle$

lemma *norm-isometry-compose'*: $\langle norm\ (A\ o_{CL}\ U) = norm\ A \rangle$ **if** $\langle isometry\ (U*) \rangle$
 $\langle proof \rangle$

lemma *unitary-nonzero[simp]*: $\langle \neg\ unitary\ (0 :: 'a :: \{chilbert-space, not-singleton\} \Rightarrow_{CL}\ -) \rangle$
 $\langle proof \rangle$

lemma *isometry-inj*:
assumes $\langle isometry\ U \rangle$
shows $\langle inj-on\ U\ X \rangle$
 $\langle proof \rangle$

lemma *unitary-inj*:
assumes $\langle unitary\ U \rangle$
shows $\langle inj-on\ U\ X \rangle$
 $\langle proof \rangle$

lemma *unitary-adj-inv*: $\langle unitary\ U \implies cblinfun-apply\ (U*) = inv\ (cblinfun-apply\ U) \rangle$
 $\langle proof \rangle$

lemma *isometry-cinner-both-sides*:
assumes $\langle isometry\ U \rangle$
shows $\langle cinner\ (U\ x)\ (U\ y) = cinner\ x\ y \rangle$
 $\langle proof \rangle$

lemma *isometry-image-is-ortho-set*:
assumes $\langle is-ortho-set\ A \rangle$
assumes $\langle isometry\ U \rangle$
shows $\langle is-ortho-set\ (U\ ` A) \rangle$
 $\langle proof \rangle$

13.7 Product spaces

lift-definition $cblinfun-left :: \langle 'a::complex-normed-vector \Rightarrow_{CL} ('a \times 'b::complex-normed-vector) \rangle$
is $\langle (\lambda x. (x, 0)) \rangle$
 $\langle proof \rangle$

lift-definition $cblinfun-right :: \langle 'b::complex-normed-vector \Rightarrow_{CL} ('a::complex-normed-vector \times 'b) \rangle$
is $\langle (\lambda x. (0, x)) \rangle$
 $\langle proof \rangle$

lemma $isometry-cblinfun-left[simp]: \langle isometry\ cblinfun-left \rangle$
 $\langle proof \rangle$

lemma $isometry-cblinfun-right[simp]: \langle isometry\ cblinfun-right \rangle$
 $\langle proof \rangle$

lemma $cblinfun-left-right-ortho[simp]: \langle cblinfun-left *_{CL} cblinfun-right = 0 \rangle$
 $\langle proof \rangle$

lemma $cblinfun-right-left-ortho[simp]: \langle cblinfun-right *_{CL} cblinfun-left = 0 \rangle$
 $\langle proof \rangle$

lemma $cblinfun-left-apply[simp]: \langle cblinfun-left *_{\mathcal{V}} \psi = (\psi, 0) \rangle$
 $\langle proof \rangle$

lemma $cblinfun-left-adj-apply[simp]: \langle cblinfun-left *_{\mathcal{V}} \psi = fst\ \psi \rangle$
 $\langle proof \rangle$

lemma $cblinfun-right-apply[simp]: \langle cblinfun-right *_{\mathcal{V}} \psi = (0, \psi) \rangle$
 $\langle proof \rangle$

lemma $cblinfun-right-adj-apply[simp]: \langle cblinfun-right *_{\mathcal{V}} \psi = snd\ \psi \rangle$
 $\langle proof \rangle$

lift-definition $ccsubspace-Times :: \langle 'a::complex-normed-vector\ ccsubspace \Rightarrow 'b::complex-normed-vector\ ccsubspace \Rightarrow ('a \times 'b)\ ccsubspace \rangle$ **is**
 $\langle \lambda S\ T. S \times T \rangle$
 $\langle proof \rangle$

lemma $ccspan-Times: \langle ccspan\ (S \times T) = ccsubspace-Times\ (ccspan\ S)\ (ccspan\ T) \rangle$ **if** $\langle 0 \in S \rangle$ **and** $\langle 0 \in T \rangle$
 $\langle proof \rangle$

lemma $ccspan-Times-sing1: \langle ccspan\ (\{0::'a::complex-normed-vector\} \times B) = ccsubspace-Times\ 0\ (ccspan\ B) \rangle$
 $\langle proof \rangle$

lemma $ccspan-Times-sing2: \langle ccspan\ (B \times \{0::'a::complex-normed-vector\}) = ccsubspace-Times\ (ccspan\ B)\ 0 \rangle$
 $\langle proof \rangle$

lemma *ccsubspace-Times-sup*: $\langle \text{sup } (\text{ccsubspace-Times } A \ B) \ (\text{ccsubspace-Times } C \ D) = \text{ccsubspace-Times } (\text{sup } A \ C) \ (\text{sup } B \ D) \rangle$
 $\langle \text{proof} \rangle$

lemma *ccsubspace-Times-top-top[simp]*: $\langle \text{ccsubspace-Times } \text{top } \text{top} = \text{top} \rangle$
 $\langle \text{proof} \rangle$

lemma *is-ortho-set-prod*:
assumes $\langle \text{is-ortho-set } B \rangle \ \langle \text{is-ortho-set } B' \rangle$
shows $\langle \text{is-ortho-set } ((B \times \{0\}) \cup (\{0\} \times B')) \rangle$
 $\langle \text{proof} \rangle$

lemma *ccsubspace-Times-ccspan*:
assumes $\langle \text{ccspan } B = S \rangle$ **and** $\langle \text{ccspan } B' = S' \rangle$
shows $\langle \text{ccspan } ((B \times \{0\}) \cup (\{0\} \times B')) = \text{ccsubspace-Times } S \ S' \rangle$
 $\langle \text{proof} \rangle$

lemma *is-onb-prod*:
assumes $\langle \text{is-onb } B \rangle \ \langle \text{is-onb } B' \rangle$
shows $\langle \text{is-onb } ((B \times \{0\}) \cup (\{0\} \times B')) \rangle$
 $\langle \text{proof} \rangle$

13.8 Images

The following definition defines the image of a closed subspace S under a bounded operator A . We do not define that image as the image of A seen as a function ($A \restriction S$) but as the topological closure of that image. This is because $A \restriction S$ might in general not be closed.

For example, if e_i ($i \in \mathbb{N}$) form an orthonormal basis, and A maps e_i to e_i/i , then all e_i are in $A \restriction S$, so the closure of $A \restriction S$ is the whole space. However, $\sum_i e_i/i$ is not in $A \restriction S$ because its preimage would have to be $\sum_i e_i$ which does not converge. So $A \restriction S$ does not contain the whole space, hence it is not closed.

lift-definition *cblinfun-image* :: $\langle 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'b::\text{complex-normed-vector} \Rightarrow 'a \ \text{ccsubspace} \Rightarrow 'b \ \text{ccsubspace} \rangle$ (**infixr** $\langle *_S \rangle \ 70$)
is $\lambda A \ S. \text{closure } (A \restriction S)$
 $\langle \text{proof} \rangle$

lemma *cblinfun-image-mono*:
assumes $a1: S \leq T$
shows $A *_S S \leq A *_S T$
 $\langle \text{proof} \rangle$

lemma *cblinfun-image-0[simp]*:
shows $U *_S 0 = 0$
thm *zero-ccsubspace-def*
 $\langle \text{proof} \rangle$

lemma *cblinfun-image-bot[simp]*: $U *_S \text{bot} = \text{bot}$
 ⟨proof⟩

lemma *cblinfun-image-sup[simp]*:
 fixes $A B :: \langle 'a :: \text{hilbert-space ccspace} \rangle$ and $U :: 'a \Rightarrow_{CL} 'b :: \text{hilbert-space}$
 shows $\langle U *_S (\text{sup } A B) = \text{sup } (U *_S A) (U *_S B) \rangle$
 ⟨proof⟩

lemma *scaleC-cblinfun-image[simp]*:
 fixes $A :: \langle 'a :: \text{hilbert-space} \Rightarrow_{CL} 'b :: \text{hilbert-space} \rangle$
 and $S :: \langle 'a \text{ ccspace} \rangle$ and $\alpha :: \text{complex}$
 shows $\langle (\alpha *_C A) *_S S = \alpha *_C (A *_S S) \rangle$
 ⟨proof⟩

lemma *cblinfun-image-id[simp]*:
 $\text{id-cblinfun} *_S \psi = \psi$
 ⟨proof⟩

lemma *cblinfun-compose-image*:
 $\langle (A \circ_{CL} B) *_S S = A *_S (B *_S S) \rangle$
 ⟨proof⟩

lemmas *cblinfun-assoc-left = cblinfun-compose-assoc[symmetric] cblinfun-compose-image[symmetric]*
 $\text{add.assoc}[\text{where } ?'a = 'a :: \text{hilbert-space} \Rightarrow_{CL} 'b :: \text{hilbert-space}, \text{symmetric}]$
lemmas *cblinfun-assoc-right = cblinfun-compose-assoc cblinfun-compose-image*
 $\text{add.assoc}[\text{where } ?'a = 'a :: \text{hilbert-space} \Rightarrow_{CL} 'b :: \text{hilbert-space}]$

lemma *cblinfun-image-INF-leq[simp]*:
 fixes $U :: 'b :: \text{complex-normed-vector} \Rightarrow_{CL} 'c :: \text{complex-normed-vector}$
 and $V :: 'a \Rightarrow 'b \text{ ccspace}$
 shows $\langle U *_S (\text{INF } i \in X. V i) \leq (\text{INF } i \in X. U *_S (V i)) \rangle$
 ⟨proof⟩

lemma *isometry-cblinfun-image-inf-distrib'*:
 fixes $U :: \langle 'a :: \text{complex-normed-vector} \Rightarrow_{CL} 'b :: \text{cbanach} \rangle$ and $B C :: 'a \text{ ccspace}$
 shows $U *_S (\text{inf } B C) \leq \text{inf } (U *_S B) (U *_S C)$
 ⟨proof⟩

lemma *cblinfun-image-eq*:
 fixes $S :: 'a :: \text{cbanach ccspace}$
 and $A B :: 'a :: \text{cbanach} \Rightarrow_{CL} 'b :: \text{cbanach}$
 assumes $\bigwedge x. x \in G \implies A *_V x = B *_V x$ and $\text{ccspan } G \geq S$
 shows $A *_S S = B *_S S$
 ⟨proof⟩

lemma *cblinfun-fixes-range*:
 assumes $A \circ_{CL} B = B$ and $\psi \in \text{space-as-set } (B *_S \text{top})$
 shows $A *_V \psi = \psi$

$\langle \text{proof} \rangle$

lemma *zero-cblinfun-image[simp]*: $0 *_S S = (0::\text{ccsubspace})$
 $\langle \text{proof} \rangle$

lemma *cblinfun-image-INF-eq-general*:
fixes $V :: 'a \Rightarrow 'b::\text{chilbert-space ccsubspace}$
and $U :: 'b \Rightarrow_{CL} 'c::\text{chilbert-space}$
and $Uinv :: 'c \Rightarrow_{CL} 'b$
assumes $Uinv U Uinv: Uinv \circ_{CL} U \circ_{CL} Uinv = Uinv$ **and** $U Uinv U: U \circ_{CL} Uinv \circ_{CL} U = U$
— Meaning: $Uinv$ is a Pseudoinverse of U
and $V: \bigwedge i. V i \leq Uinv *_S top$
and $\langle X \neq \{\} \rangle$
shows $U *_S (\text{INF } i \in X. V i) = (\text{INF } i \in X. U *_S V i)$
 $\langle \text{proof} \rangle$

lemma *unitary-range[simp]*:
assumes *unitary* U
shows $U *_S top = top$
 $\langle \text{proof} \rangle$

lemma *range-adjoint-isometry*:
assumes *isometry* U
shows $U^* *_S top = top$
 $\langle \text{proof} \rangle$

lemma *cblinfun-image-INF-eq[simp]*:
fixes $V :: 'a \Rightarrow 'b::\text{chilbert-space ccsubspace}$
and $U :: 'b \Rightarrow_{CL} 'c::\text{chilbert-space}$
assumes $\langle \text{isometry } U \rangle \langle X \neq \{\} \rangle$
shows $U *_S (\text{INF } i \in X. V i) = (\text{INF } i \in X. U *_S V i)$
 $\langle \text{proof} \rangle$

lemma *isometry-cblinfun-image-inf-distrib[simp]*:
fixes $U::\langle 'a::\text{chilbert-space} \Rightarrow_{CL} 'b::\text{chilbert-space} \rangle$
and $X Y::'a \text{ ccsubspace}$
assumes *isometry* U
shows $U *_S (\text{inf } X Y) = \text{inf } (U *_S X) (U *_S Y)$
 $\langle \text{proof} \rangle$

lemma *cblinfun-image-ccspan*:
shows $A *_S \text{ccspan } G = \text{ccspan } ((*_V) A \text{ ' } G)$
 $\langle \text{proof} \rangle$

lemma *cblinfun-apply-in-image[simp]*: $A *_V \psi \in \text{space-as-set } (A *_S \top)$
 $\langle \text{proof} \rangle$

lemma *cblinfun-plus-image-distr*:

$\langle (A + B) *_S S \leq A *_S S \sqcup B *_S S \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-sum-image-distr*:

$\langle (\sum_{i \in I}. A \ i) *_S S \leq (SUP \ i \in I. A \ i *_S S) \rangle$

$\langle \text{proof} \rangle$

lemma *space-as-set-image-commute*:

assumes $UV: \langle U \ o_{CL} \ V = id\text{-}cblinfun \rangle$ **and** $VU: \langle V \ o_{CL} \ U = id\text{-}cblinfun \rangle$

shows $\langle (*_V) \ U \text{ ' } space\text{-}as\text{-}set \ T = space\text{-}as\text{-}set \ (U *_S T) \rangle$

$\langle \text{proof} \rangle$

lemma *right-total-rel-ccsubspace*:

fixes $R :: \langle 'a::complex\text{-}normed\text{-}vector \Rightarrow 'b::complex\text{-}normed\text{-}vector \Rightarrow bool \rangle$

assumes $UV: \langle U \ o_{CL} \ V = id\text{-}cblinfun \rangle$

assumes $VU: \langle V \ o_{CL} \ U = id\text{-}cblinfun \rangle$

assumes $R\text{-}def: \langle \bigwedge x \ y. R \ x \ y \longleftrightarrow x = U *_V y \rangle$

shows $\langle \text{right-total} \ (\text{rel-ccsubspace} \ R) \rangle$

$\langle \text{proof} \rangle$

lemma *left-total-rel-ccsubspace*:

fixes $R :: \langle 'a::complex\text{-}normed\text{-}vector \Rightarrow 'b::complex\text{-}normed\text{-}vector \Rightarrow bool \rangle$

assumes $UV: \langle U \ o_{CL} \ V = id\text{-}cblinfun \rangle$

assumes $VU: \langle V \ o_{CL} \ U = id\text{-}cblinfun \rangle$

assumes $R\text{-}def: \langle \bigwedge x \ y. R \ x \ y \longleftrightarrow y = U *_V x \rangle$

shows $\langle \text{left-total} \ (\text{rel-ccsubspace} \ R) \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-image-bot-zero[simp]*: $\langle A *_S \text{top} = \text{bot} \longleftrightarrow A = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *surj-isometry-is-unitary*:

— This lemma is a bit stronger than its name suggests: We actually only require that the image of U is dense.

The converse is *unitary-surj*

fixes $U :: \langle 'a::chilbert\text{-}space \Rightarrow_{CL} 'b::chilbert\text{-}space \rangle$

assumes $\langle \text{isometry} \ U \rangle$

assumes $\langle U *_S \top = \top \rangle$

shows $\langle \text{unitary} \ U \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-apply-in-image'*: $A *_V \psi \in \text{space-as-set} \ (A *_S S) \text{ if } \langle \psi \in \text{space-as-set} \ S \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-image-ccspan-leqI*:

assumes $\langle \bigwedge v. v \in M \Longrightarrow A *_V v \in \text{space-as-set} \ T \rangle$

shows $\langle A *_S \text{ccspan } M \leq T \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-same-on-image*: $\langle A \psi = B \psi \rangle$ **if** *eq*: $\langle A \circ_{CL} C = B \circ_{CL} C \rangle$ **and**
mem: $\langle \psi \in \text{space-as-set } (C *_S \top) \rangle$
 $\langle \text{proof} \rangle$

lemma *lift-cblinfun-comp*:

— Utility lemma to lift a lemma of the form $a \circ_{CL} b = c$ to become a more general rewrite rule.

assumes $\langle a \circ_{CL} b = c \rangle$
shows $\langle a \circ_{CL} b = c \rangle$
and $\langle a \circ_{CL} (b \circ_{CL} d) = c \circ_{CL} d \rangle$
and $\langle a *_S (b *_S S) = c *_S S \rangle$
and $\langle a *_V (b *_V x) = c *_V x \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-image-def2*: $\langle A *_S S = \text{ccspan } ((*_V) A \text{ ‘ space-as-set } S) \rangle$
 $\langle \text{proof} \rangle$

lemma *unitary-image-onb*:

— Logically belongs in an earlier section but the proof uses results from this section.

assumes $\langle \text{is-onb } A \rangle$
assumes $\langle \text{unitary } U \rangle$
shows $\langle \text{is-onb } (U \text{ ‘ } A) \rangle$
 $\langle \text{proof} \rangle$

13.9 Sandwiches

lift-definition *sandwich* :: $\langle ('a::\text{hilbert-space} \Rightarrow_{CL} 'b::\text{complex-inner}) \Rightarrow (('a \Rightarrow_{CL} 'a) \Rightarrow_{CL} ('b \Rightarrow_{CL} 'b)) \rangle$ **is**
 $\langle \lambda(A::'a \Rightarrow_{CL} 'b) B. A \circ_{CL} B \circ_{CL} A^* \rangle$
 $\langle \text{proof} \rangle$

lemma *sandwich-0[simp]*: $\langle \text{sandwich } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *sandwich-apply*: $\langle \text{sandwich } A *_V B = A \circ_{CL} B \circ_{CL} A^* \rangle$
 $\langle \text{proof} \rangle$

lemma *sandwich-arg-compose*:

assumes $\langle \text{isometry } U \rangle$
shows $\langle \text{sandwich } U x \circ_{CL} \text{sandwich } U y = \text{sandwich } U (x \circ_{CL} y) \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-sandwich*: $\langle \text{norm } (\text{sandwich } A) = (\text{norm } A)^2 \rangle$ **for** $A :: \langle 'a::\{\text{hilbert-space}\} \Rightarrow_{CL} 'b::\{\text{complex-inner}\} \rangle$

⟨proof⟩

lemma *sandwich-apply-adj*: $\langle \text{sandwich } A (B*) = (\text{sandwich } A B)* \rangle$
 ⟨proof⟩

lemma *sandwich-id[simp]*: $\text{sandwich id-cblinfun} = \text{id-cblinfun}$
 ⟨proof⟩

lemma *sandwich-compose*: $\langle \text{sandwich } (A \circ_{CL} B) = \text{sandwich } A \circ_{CL} \text{sandwich } B \rangle$
 ⟨proof⟩

lemma *sandwich-scaleC-left*: $\langle \text{sandwich } (c *_C e) = (c \text{mod } c)^{\wedge 2} *_C \text{sandwich } e \rangle$
 ⟨proof⟩

lemma *sandwich-scaleR-left*: $\langle \text{sandwich } (r *_R e) = r^{\wedge 2} *_R \text{sandwich } e \rangle$
 ⟨proof⟩

lemma *inj-sandwich-isometry*: $\langle \text{inj } (\text{sandwich } U) \rangle$ **if** [simp]: $\langle \text{isometry } U \rangle$ **for** U
 $:: \langle 'a::\text{chilbert-space} \Rightarrow_{CL} 'b::\text{chilbert-space} \rangle$
 ⟨proof⟩

lemma *sandwich-isometry-id*: $\langle \text{isometry } (U*) \implies \text{sandwich } U \text{id-cblinfun} = \text{id-cblinfun} \rangle$
 ⟨proof⟩

13.10 Projectors

lift-definition *Proj* :: $('a::\text{chilbert-space}) \text{ccsubspace} \Rightarrow 'a \Rightarrow_{CL} 'a$
is $\langle \text{projection} \rangle$
 ⟨proof⟩

lemma *Proj-range[simp]*: $\text{Proj } S *_S \text{top} = S$
 ⟨proof⟩

lemma *adj-Proj*: $\langle (\text{Proj } M)* = \text{Proj } M \rangle$
 ⟨proof⟩

lemma *Proj-idempotent[simp]*: $\langle \text{Proj } M \circ_{CL} \text{Proj } M = \text{Proj } M \rangle$
 ⟨proof⟩

lift-definition *is-Proj* :: $\langle 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'a \Rightarrow \text{bool} \rangle$ **is**
 $\langle \lambda P. \exists M. \text{is-projection-on } P M \rangle$ ⟨proof⟩

lemma *is-Proj-id[simp]*: $\langle \text{is-Proj id-cblinfun} \rangle$
 ⟨proof⟩

lemma *Proj-top[simp]*: $\langle \text{Proj } \top = \text{id-cblinfun} \rangle$
 ⟨proof⟩

lemma *Proj-on-own-range'*:

fixes $P :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'a \rangle$
assumes $\langle P \circ_{CL} P = P \rangle$ **and** $\langle P = P* \rangle$
shows $\langle \text{Proj } (P *_S \text{top}) = P \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-range-closed*:
assumes *is-Proj* P
shows *closed* (*range* (*cblinfun-apply* P))
 $\langle \text{proof} \rangle$

lemma *Proj-is-Proj[simp]*:
fixes $M :: \langle 'a :: \text{chilbert-space} \text{ cccsubspace} \rangle$
shows $\langle \text{is-Proj } (\text{Proj } M) \rangle$
 $\langle \text{proof} \rangle$

lemma *is-Proj-algebraic*:
fixes $P :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'a \rangle$
shows $\langle \text{is-Proj } P \iff P \circ_{CL} P = P \wedge P = P* \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-on-own-range*:
fixes $P :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'a \rangle$
assumes $\langle \text{is-Proj } P \rangle$
shows $\langle \text{Proj } (P *_S \text{top}) = P \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-image-leq*: $(\text{Proj } S) *_S A \leq S$
 $\langle \text{proof} \rangle$

lemma *Proj-sandwich*:
fixes $A :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \text{chilbert-space} \rangle$
assumes *isometry* A
shows *sandwich* $A *_V \text{Proj } S = \text{Proj } (A *_S S)$
 $\langle \text{proof} \rangle$

lemma *Proj-orthog-ccspan-union*:
assumes $\bigwedge x y. x \in X \implies y \in Y \implies \text{is-orthogonal } x y$
shows $\langle \text{Proj } (\text{ccspan } (X \cup Y)) = \text{Proj } (\text{ccspan } X) + \text{Proj } (\text{ccspan } Y) \rangle$
 $\langle \text{proof} \rangle$

abbreviation *proj* :: $'a :: \text{chilbert-space} \Rightarrow 'a \Rightarrow_{CL} 'a$ **where** *proj* $\psi \equiv \text{Proj } (\text{ccspan } \{\psi\})$

lemma *proj-0[simp]*: $\langle \text{proj } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *ccsubspace-supI-via-Proj*:
fixes $A B C :: \langle 'a :: \text{chilbert-space} \text{ cccsubspace} \rangle$
assumes *a1*: $\langle \text{Proj } (\neg C) *_S A \leq B \rangle$

shows $A \leq B \sqcup C$
 $\langle \text{proof} \rangle$

lemma *is-Proj-idempotent*:
assumes *is-Proj* P
shows $P \circ_{CL} P = P$
 $\langle \text{proof} \rangle$

lemma *is-proj-selfadj*:
assumes *is-Proj* P
shows $P^* = P$
 $\langle \text{proof} \rangle$

lemma *is-Proj-I*:
assumes $P \circ_{CL} P = P$ **and** $P^* = P$
shows *is-Proj* P
 $\langle \text{proof} \rangle$

lemma *is-Proj-0[simp]*: *is-Proj* 0
 $\langle \text{proof} \rangle$

lemma *is-Proj-complement[simp]*:
fixes $P :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'a \rangle$
assumes $a1: \text{is-Proj } P$
shows *is-Proj* $(\text{id-cblinfun} - P)$
 $\langle \text{proof} \rangle$

lemma *Proj-bot[simp]*: *Proj* $\text{bot} = 0$
 $\langle \text{proof} \rangle$

lemma *Proj-ortho-compl*:
 $\text{Proj } (- X) = \text{id-cblinfun} - \text{Proj } X$
 $\langle \text{proof} \rangle$

lemma *Proj-inj*:
assumes $\text{Proj } X = \text{Proj } Y$
shows $X = Y$
 $\langle \text{proof} \rangle$

lemma *norm-Proj-leq1*: $\langle \text{norm } (\text{Proj } M) \leq 1 \rangle$ **for** $M :: \langle 'a :: \text{chilbert-space ccspace} \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-orthog-ccspan-insert*:
assumes $\bigwedge y. y \in Y \implies \text{is-orthogonal } x y$
shows $\langle \text{Proj } (\text{ccspan } (\text{insert } x Y)) = \text{proj } x + \text{Proj } (\text{ccspan } Y) \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-fixes-image*: $\langle \text{Proj } S *_V \psi = \psi \rangle$ **if** $\langle \psi \in \text{space-as-set } S \rangle$

$\langle \text{proof} \rangle$

lemma *norm-is-Proj*: $\langle \text{norm } P \leq 1 \rangle$ **if** $\langle \text{is-Proj } P \rangle$ **for** $P :: \langle 'a :: \text{hilbert-space} \Rightarrow_{CL} 'a \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-sup*: $\langle \text{orthogonal-spaces } S \ T \implies \text{Proj } (\text{sup } S \ T) = \text{Proj } S + \text{Proj } T \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-sum-spaces*:
assumes $\langle \text{finite } X \rangle$
assumes $\langle \bigwedge x \ y. x \in X \implies y \in X \implies x \neq y \implies \text{orthogonal-spaces } (J \ x) \ (J \ y) \rangle$
shows $\langle \text{Proj } (\sum x \in X. J \ x) = (\sum x \in X. \text{Proj } (J \ x)) \rangle$
 $\langle \text{proof} \rangle$

lemma *is-Proj-reduces-norm*:
fixes $P :: \langle 'a :: \text{complex-inner} \Rightarrow_{CL} 'a \rangle$
assumes $\langle \text{is-Proj } P \rangle$
shows $\langle \text{norm } (P *_{\mathcal{V}} \psi) \leq \text{norm } \psi \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-Proj-apply*: $\langle \text{norm } (\text{Proj } T *_{\mathcal{V}} \psi) = \text{norm } \psi \longleftrightarrow \psi \in \text{space-as-set } T \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-Proj-apply-1*: $\langle \text{norm } \psi = 1 \implies \text{norm } (\text{Proj } T *_{\mathcal{V}} \psi) = 1 \longleftrightarrow \psi \in \text{space-as-set } T \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-is-Proj-nonzero*: $\langle \text{norm } P = 1 \rangle$ **if** $\langle \text{is-Proj } P \rangle$ **and** $\langle P \neq 0 \rangle$ **for** $P :: \langle 'a :: \text{hilbert-space} \Rightarrow_{CL} 'a \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-compose-cancell*:
assumes $\langle A *_{\mathcal{S}} \top \leq S \rangle$
shows $\langle \text{Proj } S \ o_{CL} \ A = A \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-o-Proj-subspace-right*:
assumes $\langle A \geq B \rangle$
shows $\langle \text{Proj } A \ o_{CL} \ \text{Proj } B = \text{Proj } B \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-o-Proj-subspace-left*:
assumes $\langle A \leq B \rangle$
shows $\langle \text{Proj } A \ o_{CL} \ \text{Proj } B = \text{Proj } A \rangle$
 $\langle \text{proof} \rangle$

lemma *space-as-setI-via-Proj*:
assumes $\langle \text{Proj } M *_{\mathcal{V}} x = x \rangle$
shows $\langle x \in \text{space-as-set } M \rangle$
 $\langle \text{proof} \rangle$

lemma *unitary-image-ortho-compl*:
— Logically, this lemma belongs in an earlier section but its proof uses projectors.
fixes $U :: \langle 'a::\text{hilbert-space} \Rightarrow_{CL} 'b::\text{hilbert-space} \rangle$
assumes $[simp]: \langle \text{unitary } U \rangle$
shows $\langle U *_S (- A) = - (U *_S A) \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-on-image [simp]*: $\langle \text{Proj } S *_S S = S \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-nearest*:
assumes $\langle x \in \text{space-as-set } S \rangle$
shows $\langle \text{dist } (\text{Proj } S \ m) \ m \leq \text{dist } x \ m \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-0-compl*: $\langle \text{Proj } S \ x = 0 \rangle$ **if** $\langle x \in \text{space-as-set } (-S) \rangle$
 $\langle \text{proof} \rangle$

13.11 Kernel / eigenspaces

lift-definition *kernel* :: $'a::\text{complex-normed-vector} \Rightarrow_{CL} 'b::\text{complex-normed-vector}$
 $\Rightarrow 'a \text{ ccspace}$
is $\lambda f. f - \{0\}$
 $\langle \text{proof} \rangle$

definition *eigenspace* :: $\text{complex} \Rightarrow 'a::\text{complex-normed-vector} \Rightarrow_{CL} 'a \Rightarrow 'a \text{ ccspace}$
where
 $\text{eigenspace } a \ A = \text{kernel } (A - a *_C \text{id-cblinfun})$

lemma *kernel-scaleC[simp]*: $a \neq 0 \implies \text{kernel } (a *_C A) = \text{kernel } A$
for $a :: \text{complex}$ **and** $A :: (-, -) \text{ cblinfun}$
 $\langle \text{proof} \rangle$

lemma *kernel-0[simp]*: $\text{kernel } 0 = \text{top}$
 $\langle \text{proof} \rangle$

lemma *kernel-id[simp]*: $\text{kernel } \text{id-cblinfun} = 0$
 $\langle \text{proof} \rangle$

lemma *eigenspace-scaleC[simp]*:
assumes $a1: a \neq 0$
shows $\text{eigenspace } b \ (a *_C A) = \text{eigenspace } (b/a) \ A$
 $\langle \text{proof} \rangle$

lemma *eigenspace-memberD*:
assumes $x \in \text{space-as-set } (\text{eigenspace } e \ A)$
shows $A *_{\mathcal{V}} x = e *_{\mathcal{C}} x$
 $\langle \text{proof} \rangle$

lemma *kernel-memberD*:
assumes $x \in \text{space-as-set } (\text{kernel } A)$
shows $A *_{\mathcal{V}} x = 0$
 $\langle \text{proof} \rangle$

lemma *eigenspace-memberI*:
assumes $A *_{\mathcal{V}} x = e *_{\mathcal{C}} x$
shows $x \in \text{space-as-set } (\text{eigenspace } e \ A)$
 $\langle \text{proof} \rangle$

lemma *kernel-memberI*:
assumes $A *_{\mathcal{V}} x = 0$
shows $x \in \text{space-as-set } (\text{kernel } A)$
 $\langle \text{proof} \rangle$

lemma *kernel-Proj[simp]*: $\langle \text{kernel } (\text{Proj } S) = - \ S \rangle$
 $\langle \text{proof} \rangle$

lemma *orthogonal-projectors-orthogonal-spaces*:
— Logically belongs in section "Projectors".
fixes $S \ T :: \langle 'a::\text{hilbert-space ccspace} \rangle$
shows $\langle \text{orthogonal-spaces } S \ T \longleftrightarrow \text{Proj } S \ o_{\mathcal{C}\mathcal{L}} \ \text{Proj } T = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-compose-Proj-kernel[simp]*: $\langle a \ o_{\mathcal{C}\mathcal{L}} \ \text{Proj } (\text{kernel } a) = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kernel-compl-adj-range*:
shows $\langle \text{kernel } a = - \ (a *_{\mathcal{S}} \text{top}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kernel-apply-self*: $\langle A *_{\mathcal{S}} \text{kernel } A = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *leq-kernel-iff*:
shows $\langle A \leq \text{kernel } B \longleftrightarrow B *_{\mathcal{S}} A = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-image-kernel*:
assumes $\langle C *_{\mathcal{S}} A *_{\mathcal{S}} \text{kernel } B \leq \text{kernel } B \rangle$
assumes $\langle A \ o_{\mathcal{C}\mathcal{L}} \ C = \text{id-cblinfun} \rangle$
shows $\langle A *_{\mathcal{S}} \text{kernel } B = \text{kernel } (B \ o_{\mathcal{C}\mathcal{L}} \ C) \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-image-kernel-unitary*:

assumes $\langle \text{unitary } U \rangle$

shows $\langle U *_S \text{kernel } B = \text{kernel } (B \circ_{CL} U *) \rangle$

$\langle \text{proof} \rangle$

lemma *kernel-cblinfun-compose*:

assumes $\langle \text{kernel } B = 0 \rangle$

shows $\langle \text{kernel } A = \text{kernel } (B \circ_{CL} A) \rangle$

$\langle \text{proof} \rangle$

lemma *eigenspace-0[simp]*: $\langle \text{eigenspace } 0 \ A = \text{kernel } A \rangle$

$\langle \text{proof} \rangle$

lemma *kernel-isometry*: $\langle \text{kernel } U = 0 \rangle$ **if** $\langle \text{isometry } U \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-image-eigenspace-isometry*:

assumes $[simp]: \langle \text{isometry } A \rangle$ **and** $\langle c \neq 0 \rangle$

shows $\langle A *_S \text{eigenspace } c \ B = \text{eigenspace } c \ (\text{sandwich } A \ B) \rangle$

$\langle \text{proof} \rangle$

lemma *cblinfun-image-eigenspace-unitary*:

assumes $[simp]: \langle \text{unitary } A \rangle$

shows $\langle A *_S \text{eigenspace } c \ B = \text{eigenspace } c \ (\text{sandwich } A \ B) \rangle$

$\langle \text{proof} \rangle$

lemma *kernel-member-iff*: $\langle x \in \text{space-as-set } (\text{kernel } A) \longleftrightarrow A *_V x = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kernel-square[simp]*: $\langle \text{kernel } (A * \circ_{CL} A) = \text{kernel } A \rangle$

$\langle \text{proof} \rangle$

lemma *eq-on-ccsubspaces-Sup*:

fixes $a \ b :: \langle 'a :: \text{complex-normed-vector} \Rightarrow_{CL} 'b :: \text{complex-normed-vector} \rangle$

assumes $\langle \bigwedge i \ h. i \in I \implies h \in \text{space-as-set } (X \ i) \implies a \ h = b \ h \rangle$

shows $\langle \bigwedge h. h \in \text{space-as-set } (\bigsqcup_{i \in I} X \ i) \implies a \ h = b \ h \rangle$

$\langle \text{proof} \rangle$

lemma *eq-on-ccsubspaces-sup*:

fixes $a \ b :: \langle 'a :: \text{complex-normed-vector} \Rightarrow_{CL} 'b :: \text{complex-normed-vector} \rangle$

assumes $\langle \bigwedge h \ i. h \in \text{space-as-set } S \implies a \ h = b \ h \rangle$

assumes $\langle \bigwedge h \ i. h \in \text{space-as-set } T \implies a \ h = b \ h \rangle$

shows $\langle \bigwedge h. h \in \text{space-as-set } (S \sqcup T) \implies a \ h = b \ h \rangle$

$\langle \text{proof} \rangle$

13.12 Partial isometries

definition *partial-isometry* **where**

$\langle \text{partial-isometry } A \longleftrightarrow (\forall h \in \text{space-as-set } (- \text{ kernel } A). \text{ norm } (A \ h) = \text{ norm } h) \rangle$

lemma *partial-isometryI*:

assumes $\langle \bigwedge h. h \in \text{space-as-set } (- \text{ kernel } A) \implies \text{ norm } (A \ h) = \text{ norm } h \rangle$

shows $\langle \text{partial-isometry } A \rangle$

$\langle \text{proof} \rangle$

lemma

fixes $A :: \langle 'a :: \text{hilbert-space} \Rightarrow_{CL} 'b :: \text{complex-normed-vector} \rangle$

assumes *iso*: $\langle \bigwedge \psi. \psi \in \text{space-as-set } V \implies \text{ norm } (A *_V \psi) = \text{ norm } \psi \rangle$

assumes *zero*: $\langle \bigwedge \psi. \psi \in \text{space-as-set } (- V) \implies A *_V \psi = 0 \rangle$

shows *partial-isometryI'*: $\langle \text{partial-isometry } A \rangle$

and *partial-isometry-initial*: $\langle \text{kernel } A = - V \rangle$

$\langle \text{proof} \rangle$

lemma *Proj-partial-isometry[simp]*: $\langle \text{partial-isometry } (\text{Proj } S) \rangle$

$\langle \text{proof} \rangle$

lemma *is-Proj-partial-isometry*: $\langle \text{is-Proj } P \implies \text{partial-isometry } P \rangle$ **for** $P :: \langle - ::$

$\text{hilbert-space} \Rightarrow_{CL} - \rangle$

$\langle \text{proof} \rangle$

lemma *isometry-partial-isometry*: $\langle \text{isometry } P \implies \text{partial-isometry } P \rangle$

$\langle \text{proof} \rangle$

lemma *unitary-partial-isometry*: $\langle \text{unitary } P \implies \text{partial-isometry } P \rangle$

$\langle \text{proof} \rangle$

lemma *norm-partial-isometry*:

fixes $A :: \langle 'a :: \text{hilbert-space} \Rightarrow_{CL} 'b :: \text{complex-normed-vector} \rangle$

assumes $\langle \text{partial-isometry } A \rangle$ **and** $\langle A \neq 0 \rangle$

shows $\langle \text{norm } A = 1 \rangle$

$\langle \text{proof} \rangle$

lemma *partial-isometry-adj-a-o-a*:

assumes $\langle \text{partial-isometry } a \rangle$

shows $\langle a *_{o_{CL}} a = \text{Proj } (- \text{ kernel } a) \rangle$

$\langle \text{proof} \rangle$

lemma *partial-isometry-square-proj*: $\langle \text{is-Proj } (a *_{o_{CL}} a) \rangle$ **if** $\langle \text{partial-isometry } a \rangle$

$\langle \text{proof} \rangle$

lemma *partial-isometry-adj[simp]*: $\langle \text{partial-isometry } (a *) \rangle$ **if** $\langle \text{partial-isometry } a \rangle$

for $a :: \langle 'a :: \text{hilbert-space} \Rightarrow_{CL} 'b :: \text{hilbert-space} \rangle$

$\langle \text{proof} \rangle$

13.13 Isomorphisms and inverses

definition $iso\text{-}cblinfun :: \langle 'a::complex\text{-}normed\text{-}vector, 'b::complex\text{-}normed\text{-}vector \rangle$
 $cblinfun \Rightarrow bool \rangle$ **where**
 $\langle iso\text{-}cblinfun\ A = (\exists\ B. A\ o_{CL}\ B = id\text{-}cblinfun \wedge B\ o_{CL}\ A = id\text{-}cblinfun) \rangle$

definition $\langle invertible\text{-}cblinfun\ A \longleftrightarrow (\exists\ B. B\ o_{CL}\ A = id\text{-}cblinfun) \rangle$

definition $cblinfun\text{-}inv :: \langle 'a::complex\text{-}normed\text{-}vector, 'b::complex\text{-}normed\text{-}vector \rangle$
 $cblinfun \Rightarrow ('b, 'a)\ cblinfun \rangle$ **where**
 $\langle cblinfun\text{-}inv\ A = (if\ invertible\text{-}cblinfun\ A\ then\ SOME\ B. B\ o_{CL}\ A = id\text{-}cblinfun\ else\ 0) \rangle$

lemma $cblinfun\text{-}inv\text{-}left$:
assumes $\langle invertible\text{-}cblinfun\ A \rangle$
shows $\langle cblinfun\text{-}inv\ A\ o_{CL}\ A = id\text{-}cblinfun \rangle$
 $\langle proof \rangle$

lemma $inv\text{-}cblinfun\text{-}invertible$: $\langle iso\text{-}cblinfun\ A \Longrightarrow invertible\text{-}cblinfun\ A \rangle$
 $\langle proof \rangle$

lemma $cblinfun\text{-}inv\text{-}right$:
assumes $\langle iso\text{-}cblinfun\ A \rangle$
shows $\langle A\ o_{CL}\ cblinfun\text{-}inv\ A = id\text{-}cblinfun \rangle$
 $\langle proof \rangle$

lemma $cblinfun\text{-}inv\text{-}uniq$:
assumes $A\ o_{CL}\ B = id\text{-}cblinfun$ **and** $B\ o_{CL}\ A = id\text{-}cblinfun$
shows $cblinfun\text{-}inv\ A = B$
 $\langle proof \rangle$

lemma $iso\text{-}cblinfun\text{-}unitary$: $\langle unitary\ A \Longrightarrow iso\text{-}cblinfun\ A \rangle$
 $\langle proof \rangle$

lemma $invertible\text{-}cblinfun\text{-}isometry$: $\langle isometry\ A \Longrightarrow invertible\text{-}cblinfun\ A \rangle$
 $\langle proof \rangle$

lemma $summable\text{-}cblinfun\text{-}apply\text{-}invertible$:
assumes $\langle invertible\text{-}cblinfun\ A \rangle$
shows $\langle (\lambda x. A\ *_V\ g\ x)\ summable\text{-}on\ S \longleftrightarrow g\ summable\text{-}on\ S \rangle$
 $\langle proof \rangle$

lemma $infsum\text{-}cblinfun\text{-}apply\text{-}invertible$:
assumes $\langle invertible\text{-}cblinfun\ A \rangle$
shows $\langle (\sum_{x \in S} A\ *_V\ g\ x) = A\ *_V\ (\sum_{x \in S} g\ x) \rangle$
 $\langle proof \rangle$

13.14 One-dimensional spaces

instantiation $cblinfun :: (one\text{-}dim, one\text{-}dim)\ complex\text{-}inner$ **begin**

Once we have a theory for the trace, we could instead define the Hilbert-Schmidt inner product and relax the *one-dim-sort* constraint to *(cfinite-dim, complex-normed-vector)* or similar

```

definition cinner-cblinfun (A::'a  $\Rightarrow_{CL}$  'b) (B::'a  $\Rightarrow_{CL}$  'b)
  = cnj (one-dim-iso (A *_V 1)) * one-dim-iso (B *_V 1)

instance
  <proof>
end

instantiation cblinfun :: (one-dim, one-dim) one-dim begin
lift-definition one-cblinfun :: 'a  $\Rightarrow_{CL}$  'b is one-dim-iso
  <proof>
lift-definition times-cblinfun :: 'a  $\Rightarrow_{CL}$  'b  $\Rightarrow$  'a  $\Rightarrow_{CL}$  'b  $\Rightarrow$  'a  $\Rightarrow_{CL}$  'b
  is  $\lambda f g. f \circ one-dim-iso \circ g$ 
  <proof>
lift-definition inverse-cblinfun :: 'a  $\Rightarrow_{CL}$  'b  $\Rightarrow$  'a  $\Rightarrow_{CL}$  'b is
   $\lambda f. ((*) (one-dim-iso (inverse (f 1)))) \circ one-dim-iso$ 
  <proof>
definition divide-cblinfun :: 'a  $\Rightarrow_{CL}$  'b  $\Rightarrow$  'a  $\Rightarrow_{CL}$  'b  $\Rightarrow$  'a  $\Rightarrow_{CL}$  'b where
  divide-cblinfun A B = A * inverse B
definition canonical-basis-cblinfun = [1 :: 'a  $\Rightarrow_{CL}$  'b]
definition <canonical-basis-length-cblinfun (- :: ('a  $\Rightarrow_{CL}$  'b) itself) = (1::nat)>
instance
  <proof>
end

lemma id-cblinfun-eq-1[simp]: <id-cblinfun = 1>
  <proof>

lemma one-dim-cblinfun-compose-is-times[simp]:
  fixes A :: 'a::one-dim  $\Rightarrow_{CL}$  'a and B :: 'a  $\Rightarrow_{CL}$  'a
  shows A oCL B = A * B
  <proof>

lemma scaleC-one-dim-is-times: <r *_C x = one-dim-iso r * x>
  <proof>

lemma one-comp-one-cblinfun[simp]: 1 oCL 1 = 1
  <proof>

lemma one-cblinfun-adj[simp]: 1* = 1
  <proof>

lemma scaleC-1-apply[simp]: <(x *_C 1) *_V y = x *_C y>
  <proof>

lemma cblinfun-apply-1-left[simp]: <1 *_V y = y>
  <proof>

```

lemma *of-complex-cblinfun-apply[simp]*: $\langle \text{of-complex } x *_V y = \text{one-dim-iso } (x *_C y) \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-compose-1-left[simp]*: $\langle 1 \circ_{CL} x = x \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-compose-1-right[simp]*: $\langle x \circ_{CL} 1 = x \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-iso-id-cblinfun*: $\langle \text{one-dim-iso id-cblinfun} = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-iso-id-cblinfun-eq-1*: $\langle \text{one-dim-iso id-cblinfun} = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-iso-comp-distr[simp]*: $\langle \text{one-dim-iso } (a \circ_{CL} b) = \text{one-dim-iso } a \circ_{CL} \text{one-dim-iso } b \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-iso-comp-distr-times[simp]*: $\langle \text{one-dim-iso } (a \circ_{CL} b) = \text{one-dim-iso } a * \text{one-dim-iso } b \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-iso-adjoint[simp]*: $\langle \text{one-dim-iso } (A^*) = (\text{one-dim-iso } A)^* \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-iso-adjoint-complex[simp]*: $\langle \text{one-dim-iso } (A^*) = \text{cnj } (\text{one-dim-iso } A) \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-cblinfun-compose-commute*: $\langle a \circ_{CL} b = b \circ_{CL} a \rangle$ **for** $a b :: \langle ('a :: \text{one-dim}, 'a) \text{cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *one-cblinfun-apply-one[simp]*: $\langle 1 *_V 1 = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-cblinfun-apply-is-times*:
fixes $A :: 'a :: \text{one-dim} \Rightarrow_{CL} 'b :: \text{one-dim}$ **and** $b :: 'a$
shows $A *_V b = \text{one-dim-iso } A * \text{one-dim-iso } b$
 $\langle \text{proof} \rangle$

lemma *is-onb-one-dim[simp]*: $\langle \text{norm } x = 1 \implies \text{is-onb } \{x\} \rangle$ **for** $x :: \langle - :: \text{one-dim} \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-iso-cblinfun-comp*: $\langle \text{one-dim-iso } (a \circ_{CL} b) = \text{of-complex } (\text{cinner } (a *_V 1) (b *_V 1)) \rangle$
for $a :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \text{one-dim} \rangle$ **and** $b :: \langle 'c :: \text{one-dim} \Rightarrow_{CL} 'a \rangle$

⟨proof⟩

lemma *one-dim-iso-cblinfun-apply*[simp]: ⟨one-dim-iso $\psi *_{\mathcal{V}}$ $\varphi = \text{one-dim-iso } (\text{one-dim-iso } \psi *_{\mathcal{C}} \varphi)$ ⟩
 ⟨proof⟩

13.15 Loewner order

lift-definition *heterogenous-cblinfun-id* :: ⟨'a::complex-normed-vector $\Rightarrow_{CL}$ 'b::complex-normed-vector⟩
is ⟨if bounded-clinear (heterogenous-identity :: 'a::complex-normed-vector $\Rightarrow$ 'b::complex-normed-vector)
 then heterogenous-identity else ($\lambda \cdot 0$)⟩
 ⟨proof⟩

lemma *heterogenous-cblinfun-id-def'*[simp]: heterogenous-cblinfun-id = id-cblinfun
 ⟨proof⟩

definition *heterogenous-same-type-cblinfun* (x::'a::chilbert-space itself) (y::'b::chilbert-space itself) $\longleftrightarrow$
 unitary (heterogenous-cblinfun-id :: 'a $\Rightarrow_{CL}$ 'b) $\wedge$ unitary (heterogenous-cblinfun-id
 :: 'b $\Rightarrow_{CL}$ 'a)

lemma *heterogenous-same-type-cblinfun*[simp]: ⟨heterogenous-same-type-cblinfun (x::'a::chilbert-space itself) (y::'a::chilbert-space itself)⟩
 ⟨proof⟩

instantiation *cblinfun* :: (chilbert-space, chilbert-space) ord **begin**

definition *less-eq-cblinfun-def-heterogenous*: ⟨ $A \leq B \longleftrightarrow$
 (if heterogenous-same-type-cblinfun TYPE('a) TYPE('b) then
 $\forall \psi :: 'b. \psi \cdot_{\mathcal{C}} ((B-A) *_{\mathcal{V}} \text{heterogenous-cblinfun-id} *_{\mathcal{V}} \psi) \geq 0$ else ($A=B$))⟩

definition ⟨ $(A :: 'a \Rightarrow_{CL} 'b) < B \longleftrightarrow A \leq B \wedge \neg B \leq A$ ⟩

instance(proof)

end

lemma *less-eq-cblinfun-def*: ⟨ $A \leq B \longleftrightarrow$
 $(\forall \psi. \psi \cdot_{\mathcal{C}} (A *_{\mathcal{V}} \psi) \leq \psi \cdot_{\mathcal{C}} (B *_{\mathcal{V}} \psi))$ ⟩
 ⟨proof⟩

instantiation *cblinfun* :: (chilbert-space, chilbert-space) ordered-complex-vector **begin**
instance
 ⟨proof⟩
end

instance *cblinfun* :: (chilbert-space, chilbert-space) ordered-comm-monoid-add
 ⟨proof⟩

lemma *positive-id-cblinfun*[simp]: id-cblinfun ≥ 0
 ⟨proof⟩

lemma *positive-selfadjointI*: $\langle \text{selfadjoint } A \rangle$ **if** $\langle A \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-leI*:
assumes $\langle \bigwedge x. \text{norm } x = 1 \implies x \cdot_C (A *_V x) \leq x \cdot_C (B *_V x) \rangle$
shows $\langle A \leq B \rangle$
 $\langle \text{proof} \rangle$

lemma *positive-cblinfunI*: $\langle A \geq 0 \rangle$ **if** $\langle \bigwedge x. \text{norm } x = 1 \implies \text{cinner } x (A *_V x) \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *less-eq-scaled-id-norm*:
assumes $\langle \text{norm } A \leq c \rangle$ **and** $\langle \text{selfadjoint } A \rangle$
shows $\langle A \leq c *_R \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *positive-cblinfun-squareI*: $\langle A = B *_O o_{CL} B \implies A \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-loewner-order*: $\langle A \geq B \longleftrightarrow \text{one-dim-iso } A \geq (\text{one-dim-iso } B :: \text{complex}) \rangle$ **for** $A \ B :: \langle 'a \Rightarrow_{CL} 'a :: \{\text{hilbert-space, one-dim}\} \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-positive*: $\langle A \geq 0 \longleftrightarrow \text{one-dim-iso } A \geq (0 :: \text{complex}) \rangle$ **for** $A :: \langle 'a \Rightarrow_{CL} 'a :: \{\text{hilbert-space, one-dim}\} \rangle$
 $\langle \text{proof} \rangle$

lemma *op-square-nondegenerate*: $\langle a = 0 \rangle$ **if** $\langle a *_O o_{CL} a = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *comparable-selfadjoint*:
assumes $\langle a \leq b \rangle$
assumes $\langle \text{selfadjoint } a \rangle$
shows $\langle \text{selfadjoint } b \rangle$
 $\langle \text{proof} \rangle$

lemma *comparable-selfadjoint'*:
assumes $\langle a \leq b \rangle$
assumes $\langle \text{selfadjoint } b \rangle$
shows $\langle \text{selfadjoint } a \rangle$
 $\langle \text{proof} \rangle$

lemma *is-Proj-leq-id*: $\langle \text{is-Proj } P \implies P \leq \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *Proj-mono*: $\langle \text{Proj } S \leq \text{Proj } T \longleftrightarrow S \leq T \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-mono-neutral-cblinfun*:
fixes $f :: 'a \Rightarrow ('b::\text{hilbert-space} \Rightarrow_{CL} 'b)$
assumes $\langle f \text{ has-sum } a \rangle$ **and** $\langle g \text{ has-sum } b \rangle$ B
assumes $\langle \bigwedge x. x \in A \cap B \implies f\ x \leq g\ x \rangle$
assumes $\langle \bigwedge x. x \in A - B \implies f\ x \leq 0 \rangle$
assumes $\langle \bigwedge x. x \in B - A \implies g\ x \geq 0 \rangle$
shows $a \leq b$
 $\langle \text{proof} \rangle$

lemma *sums-mono-cblinfun*:
fixes $f :: \text{nat} \Rightarrow ('b::\text{hilbert-space} \Rightarrow_{CL} 'b)$
assumes $\langle f \text{ sums } a \rangle$ **and** $\langle g \text{ sums } b \rangle$
assumes $\langle \bigwedge n. f\ n \leq g\ n \rangle$
shows $a \leq b$
 $\langle \text{proof} \rangle$

13.16 Embedding vectors to operators

lift-definition *vector-to-cblinfun* :: $'a::\text{complex-normed-vector} \Rightarrow 'b::\text{one-dim} \Rightarrow_{CL}$
 $'a \rangle$ **is**
 $\langle \lambda \psi. \varphi. \text{one-dim-iso } \varphi *_C \psi \rangle$
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-apply[simp]*: $\langle \text{vector-to-cblinfun } \psi *_V \varphi = \text{one-dim-iso } \psi *_C \varphi \rangle$
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-cblinfun-compose[simp]*:
 $A \circ_{CL} (\text{vector-to-cblinfun } \psi) = \text{vector-to-cblinfun } (A *_V \psi)$
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-add*: $\langle \text{vector-to-cblinfun } (x + y) = \text{vector-to-cblinfun } x + \text{vector-to-cblinfun } y \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-vector-to-cblinfun[simp]*: $\text{norm } (\text{vector-to-cblinfun } x) = \text{norm } x$
 $\langle \text{proof} \rangle$

lemma *bounded-clinear-vector-to-cblinfun[bounded-clinear]*: *bounded-clinear* *vector-to-cblinfun*
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-scaleC[simp]*:
 $\text{vector-to-cblinfun } (a *_C \psi) = a *_C \text{vector-to-cblinfun } \psi$ **for** $a::\text{complex}$
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-apply-one-dim[simp]*:
shows $\text{vector-to-cblinfun } \varphi *_V \gamma = \text{one-dim-iso } \gamma *_C \varphi$
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-one-dim-iso*[simp]: $\langle \text{vector-to-cblinfun} = \text{one-dim-iso} \rangle$
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-adj-apply*[simp]:
shows $\text{vector-to-cblinfun } \psi * *_V \varphi = \text{of-complex } (\text{cinner } \psi \varphi)$
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-comp-one*[simp]:
 $(\text{vector-to-cblinfun } s :: 'a::\text{one-dim} \Rightarrow_{CL} -) \circ_{CL} 1$
 $= (\text{vector-to-cblinfun } s :: 'b::\text{one-dim} \Rightarrow_{CL} -)$
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-0*[simp]: $\text{vector-to-cblinfun } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *image-vector-to-cblinfun*[simp]: $\text{vector-to-cblinfun } x *_S \top = \text{ccspan } \{x\}$
— Not that the general case $\text{vector-to-cblinfun } x *_S S$ can be handled by using
that $S = \top$ or $S = \perp$ by *one-dim-ccsubspace-all-or-nothing*
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-adj-comp-vector-to-cblinfun*[simp]:
shows $\text{vector-to-cblinfun } \psi * \circ_{CL} \text{vector-to-cblinfun } \varphi = \text{cinner } \psi \varphi *_C \text{id-cblinfun}$
 $\langle \text{proof} \rangle$

lemma *isometry-vector-to-cblinfun*[simp]:
assumes $\text{norm } x = 1$
shows $\text{isometry } (\text{vector-to-cblinfun } x)$
 $\langle \text{proof} \rangle$

lemma *image-vector-to-cblinfun-adj*:
assumes $\langle \psi \notin \text{space-as-set } (- S) \rangle$
shows $\langle (\text{vector-to-cblinfun } \psi) * *_S S = \top \rangle$
 $\langle \text{proof} \rangle$

lemma *image-vector-to-cblinfun-adj'*:
assumes $\langle \psi \neq 0 \rangle$
shows $\langle (\text{vector-to-cblinfun } \psi) * *_S \top = \top \rangle$
 $\langle \text{proof} \rangle$

lemma *vector-to-cblinfun-inj*: $\langle \text{inj-on } (\text{vector-to-cblinfun } :: 'a::\text{complex-normed-vector} \Rightarrow 'b::\text{one-dim} \Rightarrow_{CL} -) X \rangle$
 $\langle \text{proof} \rangle$

13.17 Rank-1 operators / butterflies

definition *rank1* **where** $\langle \text{rank1 } A \longleftrightarrow (\exists \psi. A *_S \top = \text{ccspan } \{\psi\}) \rangle$

— This is not the usual definition of a rank-1 operator. The usual definition is

an operator with 1-dim image. Here we define it as an operator with 0- or 1-dim image. This makes the definition simpler to use. The normal definition of rank-1 operators then corresponds to the non-zero *rank1* operators.

definition *butterfly* (*s*::'a::complex-normed-vector) (*t*::'b::hilbert-space)
 $= \text{vector-to-cblinfun } s \text{ } o_{CL} (\text{vector-to-cblinfun } t :: \text{complex} \Rightarrow_{CL} -)^*$

abbreviation *selfbutter* *s* $\equiv \text{butterfly } s \text{ } s$

lemma *butterfly-add-left*: $\langle \text{butterfly } (a + a') \text{ } b = \text{butterfly } a \text{ } b + \text{butterfly } a' \text{ } b \rangle$
 $\langle \text{proof} \rangle$

lemma *butterfly-add-right*: $\langle \text{butterfly } a \text{ } (b + b') = \text{butterfly } a \text{ } b + \text{butterfly } a \text{ } b' \rangle$
 $\langle \text{proof} \rangle$

lemma *butterfly-def-one-dim*: $\text{butterfly } s \text{ } t = (\text{vector-to-cblinfun } s :: 'c::\text{one-dim} \Rightarrow_{CL} -)$
 $o_{CL} (\text{vector-to-cblinfun } t :: 'c \Rightarrow_{CL} -)^*$
(is - = ?rhs) for *s* :: 'a::complex-normed-vector **and** *t* :: 'b::hilbert-space
 $\langle \text{proof} \rangle$

lemma *butterfly-comp-cblinfun*: $\text{butterfly } \psi \text{ } \varphi \text{ } o_{CL} \text{ } a = \text{butterfly } \psi \text{ } (a *_{\text{V}} \varphi)$
 $\langle \text{proof} \rangle$

lemma *cblinfun-comp-butterfly*: $a \text{ } o_{CL} \text{ } \text{butterfly } \psi \text{ } \varphi = \text{butterfly } (a *_{\text{V}} \psi) \text{ } \varphi$
 $\langle \text{proof} \rangle$

lemma *butterfly-apply[simp]*: $\text{butterfly } \psi \text{ } \psi' *_{\text{V}} \varphi = (\psi' \cdot_C \varphi) *_C \psi$
 $\langle \text{proof} \rangle$

lemma *butterfly-scaleC-left[simp]*: $\text{butterfly } (c *_C \psi) \text{ } \varphi = c *_C \text{butterfly } \psi \text{ } \varphi$
 $\langle \text{proof} \rangle$

lemma *butterfly-scaleC-right[simp]*: $\text{butterfly } \psi \text{ } (c *_C \varphi) = c *_C \text{butterfly } \psi \text{ } \varphi$
 $\langle \text{proof} \rangle$

lemma *butterfly-scaleR-left[simp]*: $\text{butterfly } (r *_R \psi) \text{ } \varphi = r *_C \text{butterfly } \psi \text{ } \varphi$
 $\langle \text{proof} \rangle$

lemma *butterfly-scaleR-right[simp]*: $\text{butterfly } \psi \text{ } (r *_R \varphi) = r *_C \text{butterfly } \psi \text{ } \varphi$
 $\langle \text{proof} \rangle$

lemma *butterfly-adjoint[simp]*: $(\text{butterfly } \psi \text{ } \varphi)^* = \text{butterfly } \varphi \text{ } \psi$
 $\langle \text{proof} \rangle$

lemma *butterfly-comp-butterfly[simp]*: $\text{butterfly } \psi_1 \text{ } \psi_2 \text{ } o_{CL} \text{ } \text{butterfly } \psi_3 \text{ } \psi_4 = (\psi_2 \cdot_C \psi_3) *_C \text{butterfly } \psi_1 \text{ } \psi_4$
 $\langle \text{proof} \rangle$

lemma *butterfly-0-left[simp]*: *butterfly 0 a = 0*
 $\langle \text{proof} \rangle$

lemma *butterfly-0-right[simp]*: *butterfly a 0 = 0*
 $\langle \text{proof} \rangle$

lemma *butterfly-is-rank1*:
assumes $\langle \varphi \neq 0 \rangle$
shows $\langle \text{butterfly } \psi \ \varphi *_{\mathcal{S}} \top = \text{ccspan } \{\psi\} \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-is-butterfly*:
— The restriction ψ is necessary. Consider, e.g., the space of all finite sequences (with sum-norm), and $A' f = (\sum x. f x)$. Then A' is not a butterfly.
assumes $\langle A *_{\mathcal{S}} \top = \text{ccspan } \{\psi :: \text{chilbert-space}\} \rangle$
shows $\langle \exists \varphi. A = \text{butterfly } \psi \ \varphi \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-0[simp]*: $\langle \text{rank1 } 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-iff-butterfly*: $\langle \text{rank1 } A \longleftrightarrow (\exists \psi \ \varphi. A = \text{butterfly } \psi \ \varphi) \rangle$
for $A :: \langle \text{complex-inner} \Rightarrow_{CL} \text{chilbert-space} \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-butterfly[iff]*: $\langle \text{rank1 } (\text{butterfly } x \ y) \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-butterfly*: $\text{norm } (\text{butterfly } \psi \ \varphi) = \text{norm } \psi * \text{norm } \varphi$
 $\langle \text{proof} \rangle$

lemma *bounded-sesquilinear-butterfly[bounded-sesquilinear]*: $\langle \text{bounded-sesquilinear } (\lambda(b :: \text{chilbert-space}) (a :: \text{chilbert-space}). \text{butterfly } a \ b) \rangle$
 $\langle \text{proof} \rangle$

lemma *inj-selfbutter-upto-phase*:
assumes $\text{selfbutter } x = \text{selfbutter } y$
shows $\exists c. \text{cmod } c = 1 \wedge x = c *_C y$
 $\langle \text{proof} \rangle$

lemma *butterfly-eq-proj*:
assumes $\text{norm } x = 1$
shows $\text{selfbutter } x = \text{proj } x$
 $\langle \text{proof} \rangle$

lemma *butterfly-sgn-eq-proj*:
shows $\text{selfbutter } (\text{sgn } x) = \text{proj } x$
 $\langle \text{proof} \rangle$

lemma *butterfly-is-Proj*:

$\langle \text{norm } x = 1 \implies \text{is-Proj } (\text{selfbutter } x) \rangle$

$\langle \text{proof} \rangle$

lemma *cspan-butterfly-UNIV*:

assumes $\langle \text{cspan basisA} = \text{UNIV} \rangle$

assumes $\langle \text{cspan basisB} = \text{UNIV} \rangle$

assumes $\langle \text{is-ortho-set basisB} \rangle$

assumes $\langle \bigwedge b. b \in \text{basisB} \implies \text{norm } b = 1 \rangle$

shows $\langle \text{cspan } \{ \text{butterfly } a \mid (a :: 'a :: \{ \text{complex-normed-vector} \}) (b :: 'b :: \{ \text{chilbert-space, cfinite-dim} \}).$

$a \in \text{basisA} \wedge b \in \text{basisB} \} = \text{UNIV} \rangle$

$\langle \text{proof} \rangle$

lemma *cindependent-butterfly*:

fixes $\text{basisA} :: \langle 'a :: \text{chilbert-space set} \rangle$ **and** $\text{basisB} :: \langle 'b :: \text{chilbert-space set} \rangle$

assumes $\langle \text{is-ortho-set basisA} \rangle$ $\langle \text{is-ortho-set basisB} \rangle$

assumes $\text{normA}: \langle \bigwedge a. a \in \text{basisA} \implies \text{norm } a = 1 \rangle$ **and** $\text{normB}: \langle \bigwedge b. b \in \text{basisB} \implies \text{norm } b = 1 \rangle$

shows $\langle \text{cindependent } \{ \text{butterfly } a \mid a \in \text{basisA} \wedge b \in \text{basisB} \} \rangle$

$\langle \text{proof} \rangle$

lemma *clinear-eq-butterflyI*:

fixes $F \ G :: \langle ('a :: \{ \text{chilbert-space, cfinite-dim} \} \Rightarrow_{CL} 'b :: \text{complex-inner}) \Rightarrow 'c :: \text{complex-vector} \rangle$

assumes *clinear* F **and** *clinear* G

assumes $\langle \text{cspan basisA} = \text{UNIV} \rangle$ $\langle \text{cspan basisB} = \text{UNIV} \rangle$

assumes $\langle \text{is-ortho-set basisA} \rangle$ $\langle \text{is-ortho-set basisB} \rangle$

assumes $\langle \bigwedge a \ b. a \in \text{basisA} \implies b \in \text{basisB} \implies F (\text{butterfly } a \ b) = G (\text{butterfly } a \ b) \rangle$

assumes $\langle \bigwedge b. b \in \text{basisB} \implies \text{norm } b = 1 \rangle$

shows $F = G$

$\langle \text{proof} \rangle$

lemma *sum-butterfly-is-Proj*:

assumes $\langle \text{is-ortho-set } E \rangle$

assumes $\langle \bigwedge e. e \in E \implies \text{norm } e = 1 \rangle$

shows $\langle \text{is-Proj } (\sum_{e \in E. \text{butterfly } e \ e}) \rangle$

$\langle \text{proof} \rangle$

lemma *rank1-compose-left*: $\langle \text{rank1 } (a \ o_{CL} \ b) \rangle$ **if** $\langle \text{rank1 } b \rangle$

$\langle \text{proof} \rangle$

lemma *csubspace-of-1dim-space*:

assumes $\langle S \neq \{0\} \rangle$

assumes $\langle \text{csubspace } S \rangle$

assumes $\langle S \subseteq \text{cspan } \{ \psi \} \rangle$

shows $\langle S = \text{cspan } \{ \psi \} \rangle$

$\langle \text{proof} \rangle$

lemma *subspace-of-1dim-ccspan*:

assumes $\langle S \neq 0 \rangle$

assumes $\langle S \leq \text{ccspan } \{\psi\} \rangle$
shows $\langle S = \text{ccspan } \{\psi\} \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-compose-right*: $\langle \text{rank1 } (a \circ_{CL} b) \rangle$ **if** $\langle \text{rank1 } a \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-scaleC*: $\langle \text{rank1 } (c *_C a) \rangle$ **if** $\langle \text{rank1 } a \rangle$ **and** $\langle c \neq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-scaleR[simp]*: $\langle \text{rank1 } (c *_R a) \rangle$ **if** $\langle \text{rank1 } a \rangle$ **and** $\langle c \neq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-uminus*: $\langle \text{rank1 } (-a) \rangle$ **if** $\langle \text{rank1 } a \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-uminus-iff[simp]*: $\langle \text{rank1 } (-a) \longleftrightarrow \text{rank1 } a \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-adj*: $\langle \text{rank1 } (a^*) \rangle$ **if** $\langle \text{rank1 } a \rangle$
for $a :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \text{chilbert-space} \rangle$
 $\langle \text{proof} \rangle$

lemma *rank1-adj-iff[simp]*: $\langle \text{rank1 } (a^*) \longleftrightarrow \text{rank1 } a \rangle$
for $a :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \text{chilbert-space} \rangle$
 $\langle \text{proof} \rangle$

lemma *butterflies-sum-id-finite*: $\langle \text{id-cblinfun} = (\sum x \in B. \text{selfbutter } x) \rangle$ **if** $\langle \text{is-onb } B \rangle$ **for** $B :: \langle 'a :: \{ \text{cfinite-dim}, \text{chilbert-space} \} \text{ set} \rangle$
 $\langle \text{proof} \rangle$

lemma *butterfly-sum-left*: $\langle \text{butterfly } (\sum i \in M. \psi \ i) \ \varphi = (\sum i \in M. \text{butterfly } (\psi \ i) \ \varphi) \rangle$
 $\langle \text{proof} \rangle$

lemma *butterfly-sum-right*: $\langle \text{butterfly } \psi \ (\sum i \in M. \varphi \ i) = (\sum i \in M. \text{butterfly } \psi \ (\varphi \ i)) \rangle$
 $\langle \text{proof} \rangle$

lemma *sum-butterfly-leq-id*:
assumes $\langle \text{is-ortho-set } E \rangle$
assumes $\langle \bigwedge e. e \in E \implies \text{norm } e = 1 \rangle$
shows $\langle (\sum i \in E. \text{butterfly } i \ i) \leq \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *sandwich-butterfly*: $\langle \text{sandwich } a \ (\text{butterfly } x \ y) = \text{butterfly } (a \ x) \ (a \ y) \rangle$
 $\langle \text{proof} \rangle$

13.18 Banach-Steinhaus

theorem *cbanach-steinhaus*:

fixes $F :: \langle 'c \Rightarrow 'a :: \text{cbanach} \Rightarrow_{CL} 'b :: \text{complex-normed-vector} \rangle$

assumes $\langle \bigwedge x. \exists M. \forall n. \text{norm } ((F\ n) *_{\mathcal{V}} x) \leq M \rangle$

shows $\langle \exists M. \forall n. \text{norm } (F\ n) \leq M \rangle$

$\langle \text{proof} \rangle$

13.19 Riesz-representation theorem

theorem *riesz-representation-cblinfun-existence*:

— Theorem 3.4 in [1]

fixes $f :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} \text{complex} \rangle$

shows $\langle \exists t. \forall x. f *_{\mathcal{V}} x = (t \cdot_C x) \rangle$

$\langle \text{proof} \rangle$

lemma *riesz-representation-cblinfun-unique*:

— Theorem 3.4 in [1]

fixes $f :: \langle 'a :: \text{complex-inner} \Rightarrow_{CL} \text{complex} \rangle$

assumes $\langle \bigwedge x. f *_{\mathcal{V}} x = (t \cdot_C x) \rangle$

assumes $\langle \bigwedge x. f *_{\mathcal{V}} x = (u \cdot_C x) \rangle$

shows $\langle t = u \rangle$

$\langle \text{proof} \rangle$

theorem *riesz-representation-cblinfun-norm*:

includes *norm-syntax*

fixes $f :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} \text{complex} \rangle$

assumes $\langle \bigwedge x. f *_{\mathcal{V}} x = (t \cdot_C x) \rangle$

shows $\langle \|f\| = \|t\| \rangle$

$\langle \text{proof} \rangle$

definition *the-riesz-rep* :: $\langle ('a :: \text{chilbert-space} \Rightarrow_{CL} \text{complex}) \Rightarrow 'a \rangle$ **where**

$\langle \text{the-riesz-rep } f = (\text{SOME } t. \forall x. f *_{\mathcal{V}} x = t \cdot_C x) \rangle$

lemma *the-riesz-rep[simp]*: $\langle \text{the-riesz-rep } f \cdot_C x = f *_{\mathcal{V}} x \rangle$

$\langle \text{proof} \rangle$

lemma *the-riesz-rep-unique*:

assumes $\langle \bigwedge x. f *_{\mathcal{V}} x = t \cdot_C x \rangle$

shows $\langle t = \text{the-riesz-rep } f \rangle$

$\langle \text{proof} \rangle$

lemma *the-riesz-rep-scaleC*: $\langle \text{the-riesz-rep } (c \cdot_C f) = \text{cnj } c \cdot_C \text{the-riesz-rep } f \rangle$

$\langle \text{proof} \rangle$

lemma *the-riesz-rep-add*: $\langle \text{the-riesz-rep } (f + g) = \text{the-riesz-rep } f + \text{the-riesz-rep } g \rangle$

$\langle \text{proof} \rangle$

lemma *the-riesz-rep-norm[simp]*: $\langle \text{norm } (\text{the-riesz-rep } f) = \text{norm } f \rangle$

⟨proof⟩

lemma *bounded-antilinear-the-riesz-rep*[*bounded-antilinear*]: ⟨*bounded-antilinear the-riesz-rep*⟩
 ⟨proof⟩

lift-definition *the-riesz-rep-sesqui* :: ⟨('a::complex-normed-vector ⇒ 'b::hilbert-space
 ⇒ complex) ⇒ ('a ⇒_{CL} 'b)⟩ **is**
 ⟨λp. if bounded-sesquilinear p then the-riesz-rep o CBlinfun o p else 0⟩
 ⟨proof⟩

lemma *the-riesz-rep-sesqui-apply*:
assumes ⟨bounded-sesquilinear p⟩
shows ⟨(the-riesz-rep-sesqui p *_V x) •_C y = p x y⟩
 ⟨proof⟩

13.20 Bidual

lift-definition *bidual-embedding* :: ⟨'a::complex-normed-vector ⇒_{CL} (('a ⇒_{CL} complex) ⇒_{CL} complex)⟩
is ⟨λx f. f *_V x⟩
 ⟨proof⟩

lemma *bidual-embedding-apply*[*simp*]: ⟨(bidual-embedding *_V x) *_V f = f *_V x⟩
 ⟨proof⟩

lemma *bidual-embedding-isometric*[*simp*]: ⟨norm (bidual-embedding *_V x) = norm x⟩ **for** x :: ⟨'a::complex-inner⟩
 ⟨proof⟩

lemma *norm-bidual-embedding*[*simp*]: ⟨norm (bidual-embedding :: 'a::{complex-inner, not-singleton} ⇒_{CL} -) = 1⟩
 ⟨proof⟩

lemma *isometry-bidual-embedding*[*simp*]: ⟨isometry bidual-embedding⟩
 ⟨proof⟩

lemma *bidual-embedding-surj*[*simp*]: ⟨surj (bidual-embedding :: 'a::hilbert-space ⇒_{CL} -)⟩
 ⟨proof⟩

13.21 Extension of complex bounded operators

definition *cblinfun-extension* **where**
cblinfun-extension S φ = (SOME B. ∀ x ∈ S. B *_V x = φ x)

definition *cblinfun-extension-exists* **where**
cblinfun-extension-exists S φ = (∃ B. ∀ x ∈ S. B *_V x = φ x)

lemma *cblinfun-extension-existsI*:
assumes ∧ x. x ∈ S ⇒ B *_V x = φ x

shows *cblinfun-extension-exists* $S \varphi$
 ⟨proof⟩

lemma *cblinfun-extension-exists-finite-dim*:
 fixes $\varphi :: 'a :: \{\text{complex-normed-vector}, \text{cfinite-dim}\} \Rightarrow 'b :: \text{complex-normed-vector}$
 assumes *cindependent* S
 and $\text{cspan } S = \text{UNIV}$
 shows *cblinfun-extension-exists* $S \varphi$
 ⟨proof⟩

lemma *cblinfun-extension-apply*:
 assumes *cblinfun-extension-exists* $S f$
 and $v \in S$
 shows $(\text{cblinfun-extension } S f) *_V v = f v$
 ⟨proof⟩

lemma
 fixes $f :: \langle 'a :: \text{complex-normed-vector} \Rightarrow 'b :: \text{cbanach} \rangle$
 assumes $\langle \text{csubspace } S \rangle$
 assumes $\langle \text{closure } S = \text{UNIV} \rangle$
 assumes *f-add*: $\langle \bigwedge x y. x \in S \implies y \in S \implies f (x + y) = f x + f y \rangle$
 assumes *f-scale*: $\langle \bigwedge c x y. x \in S \implies f (c *_C x) = c *_C f x \rangle$
 assumes *bounded*: $\langle \bigwedge x. x \in S \implies \text{norm } (f x) \leq B * \text{norm } x \rangle$
 shows *cblinfun-extension-exists-bounded-dense*: $\langle \text{cblinfun-extension-exists } S f \rangle$
 and *cblinfun-extension-norm-bounded-dense*: $\langle B \geq 0 \implies \text{norm } (\text{cblinfun-extension } S f) \leq B \rangle$
 ⟨proof⟩

lemma *cblinfun-extension-cong*:
 assumes $\langle \text{cspan } A = \text{cspan } B \rangle$
 assumes $\langle B \subseteq A \rangle$
 assumes *fg*: $\langle \bigwedge x. x \in B \implies f x = g x \rangle$
 assumes $\langle \text{cblinfun-extension-exists } A f \rangle$
 shows $\langle \text{cblinfun-extension } A f = \text{cblinfun-extension } B g \rangle$
 ⟨proof⟩

lemma
 fixes $f :: \langle 'a :: \text{complex-inner} \Rightarrow 'b :: \text{chilbert-space} \rangle$ and S
 assumes $\langle \text{is-ortho-set } S \rangle$ and $\langle \text{closure } (\text{cspan } S) = \text{UNIV} \rangle$
 assumes *ortho-f*: $\langle \bigwedge x y. x \in S \implies y \in S \implies x \neq y \implies \text{is-orthogonal } (f x) (f y) \rangle$
 assumes *bounded*: $\langle \bigwedge x. x \in S \implies \text{norm } (f x) \leq B * \text{norm } x \rangle$
 shows *cblinfun-extension-exists-ortho*: $\langle \text{cblinfun-extension-exists } S f \rangle$
 and *cblinfun-extension-exists-ortho-norm*: $\langle B \geq 0 \implies \text{norm } (\text{cblinfun-extension } S f) \leq B \rangle$
 ⟨proof⟩

lemma *cblinfun-extension-exists-proj*:
 fixes $f :: \langle 'a :: \text{complex-normed-vector} \Rightarrow 'b :: \text{cbanach} \rangle$

assumes $\langle csubspace\ S \rangle$
assumes $ex\text{-}P: \langle \exists P :: 'a \Rightarrow_{CL} 'a. is\text{-}Proj\ P \wedge range\ P = closure\ S \rangle$
assumes $f\text{-}add: \langle \bigwedge x\ y. x \in S \implies y \in S \implies f\ (x + y) = f\ x + f\ y \rangle$
assumes $f\text{-}scale: \langle \bigwedge c\ x\ y. x \in S \implies f\ (c *_{\mathcal{C}} x) = c *_{\mathcal{C}} f\ x \rangle$
assumes $bounded: \langle \bigwedge x. x \in S \implies norm\ (f\ x) \leq B * norm\ x \rangle$
shows $\langle cblinfun\text{-}extension\text{-}exists\ S\ f \rangle$

— We cannot give a statement about the norm. While there is an extension with norm B , there is no guarantee that $cblinfun\text{-}extension\ S\ f$ returns that specific extension since the extension is only determined on $ccspan\ S$.

$\langle proof \rangle$

lemma $cblinfun\text{-}extension\text{-}exists\text{-}hilbert$:

fixes $f :: \langle 'a::chilbert\text{-}space \Rightarrow 'b::cbanach \rangle$
assumes $\langle csubspace\ S \rangle$
assumes $f\text{-}add: \langle \bigwedge x\ y. x \in S \implies y \in S \implies f\ (x + y) = f\ x + f\ y \rangle$
assumes $f\text{-}scale: \langle \bigwedge c\ x\ y. x \in S \implies f\ (c *_{\mathcal{C}} x) = c *_{\mathcal{C}} f\ x \rangle$
assumes $bounded: \langle \bigwedge x. x \in S \implies norm\ (f\ x) \leq B * norm\ x \rangle$
shows $\langle cblinfun\text{-}extension\text{-}exists\ S\ f \rangle$

— We cannot give a statement about the norm. While there is an extension with norm B , there is no guarantee that $cblinfun\text{-}extension\ S\ f$ returns that specific extension since the extension is only determined on $ccspan\ S$.

$\langle proof \rangle$

lemma $cblinfun\text{-}extension\text{-}exists\text{-}restrict$:

assumes $\langle B \subseteq A \rangle$
assumes $\langle \bigwedge x. x \in B \implies f\ x = g\ x \rangle$
assumes $\langle cblinfun\text{-}extension\text{-}exists\ A\ f \rangle$
shows $\langle cblinfun\text{-}extension\text{-}exists\ B\ g \rangle$
 $\langle proof \rangle$

13.22 Bijections between different ONBs

Some of the theorems here logically belong into *Complex-Bounded-Operators.Complex-Inner-Product* but the proof uses some concepts from the present theory.

lemma $all\text{-}ortho\text{-}bases\text{-}same\text{-}card$:

— Follows [1], Proposition 4.14
fixes $E\ F :: \langle 'a::chilbert\text{-}space\ set \rangle$
assumes $\langle is\text{-}ortho\text{-}set\ E \rangle \langle is\text{-}ortho\text{-}set\ F \rangle \langle ccspan\ E = \top \rangle \langle ccspan\ F = \top \rangle$
shows $\langle \exists f. bij\text{-}betw\ f\ E\ F \rangle$
 $\langle proof \rangle$

lemma $all\text{-}onb\text{-}same\text{-}card$:

fixes $E\ F :: \langle 'a::chilbert\text{-}space\ set \rangle$
assumes $\langle is\text{-}onb\ E \rangle \langle is\text{-}onb\ F \rangle$
shows $\langle \exists f. bij\text{-}betw\ f\ E\ F \rangle$
 $\langle proof \rangle$

definition $bij\text{-}between\text{-}bases$ **where** $\langle bij\text{-}between\text{-}bases\ E\ F = (SOME\ f. bij\text{-}betw\ f\ E\ F) \rangle$ **for** $E\ F :: \langle 'a::chilbert\text{-}space\ set \rangle$

lemma *bij-between-bases-bij*:
fixes $E\ F :: \langle 'a::\text{chilbert-space set} \rangle$
assumes $\langle \text{is-onb } E \rangle \langle \text{is-onb } F \rangle$
shows $\langle \text{bij-betw } (\text{bij-between-bases } E\ F) \ E\ F \rangle$
 $\langle \text{proof} \rangle$

definition *unitary-between* **where** $\langle \text{unitary-between } E\ F = \text{cblinfun-extension } E$
 $(\text{bij-between-bases } E\ F) \rangle$

lemma *unitary-between-apply*:
fixes $E\ F :: \langle 'a::\text{chilbert-space set} \rangle$
assumes $\langle \text{is-onb } E \rangle \langle \text{is-onb } F \rangle \langle e \in E \rangle$
shows $\langle \text{unitary-between } E\ F *_{\mathcal{V}} e = \text{bij-between-bases } E\ F\ e \rangle$
 $\langle \text{proof} \rangle$

lemma *unitary-between-unitary*:
fixes $E\ F :: \langle 'a::\text{chilbert-space set} \rangle$
assumes $\langle \text{is-onb } E \rangle \langle \text{is-onb } F \rangle$
shows $\langle \text{unitary } (\text{unitary-between } E\ F) \rangle$
 $\langle \text{proof} \rangle$

13.23 Notation

bundle *cblinfun-syntax* **begin**
notation *cblinfun-compose* (**infixl** $\langle o_{CL} \rangle$ 67)
notation *cblinfun-apply* (**infixr** $\langle *_{\mathcal{V}} \rangle$ 70)
notation *cblinfun-image* (**infixr** $\langle *_S \rangle$ 70)
notation *adj* ($\langle \cdot \rangle$ [99] 100)
type-notation *cblinfun* ($\langle (- \Rightarrow_{CL} /-) \rangle$ [22, 21] 21)
end

unbundle *no cblinfun-syntax* **and** *no lattice-syntax*

end

14 Complex-L2 – Hilbert space of square-summable functions

theory *Complex-L2*
imports
Complex-Bounded-Linear-Function

HOL-Analysis.L2-Norm
HOL-Library.Rewrite
HOL-Analysis.Infinite-Sum
HOL-Library.Infinite-Typeclass
begin

unbundle *lattice-syntax* and *cblinfun-syntax* and *no blinfun-apply-syntax*

14.1 l2 norm of functions

definition $\langle \text{has-ell2-norm } (x :: \Rightarrow \text{complex}) \longleftrightarrow (\lambda i. (x\ i)^2) \text{ abs-summable-on } UNIV \rangle$

lemma *has-ell2-norm-bdd-above*: $\langle \text{has-ell2-norm } x \longleftrightarrow \text{bdd-above } (\text{sum } (\lambda x a. \text{norm } ((x\ a)^2))) \text{ 'Collect finite'} \rangle$
 $\langle \text{proof} \rangle$

lemma *has-ell2-norm-L2-set*: $\text{has-ell2-norm } x = \text{bdd-above } (L2\text{-set } (\text{norm } o\ x) \text{ 'Collect finite'})$
 $\langle \text{proof} \rangle$

definition $\text{ell2-norm} :: \langle ('a \Rightarrow \text{complex}) \Rightarrow \text{real} \rangle$ **where** $\langle \text{ell2-norm } f = \text{sqrt } (\sum_{\infty} x. \text{norm } (f\ x)^2) \rangle$

lemma *ell2-norm-SUP*:
assumes $\langle \text{has-ell2-norm } x \rangle$
shows $\text{ell2-norm } x = \text{sqrt } (\text{SUP } F \in \{F. \text{finite } F\}. \text{sum } (\lambda i. \text{norm } (x\ i)^2) F)$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-L2-set*:
assumes $\text{has-ell2-norm } x$
shows $\text{ell2-norm } x = (\text{SUP } F \in \{F. \text{finite } F\}. L2\text{-set } (\text{norm } o\ x) F)$
 $\langle \text{proof} \rangle$

lemma *has-ell2-norm-finite[simp]*: $\text{has-ell2-norm } (f :: 'a :: \text{finite} \Rightarrow -)$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-finite*:
 $\text{ell2-norm } (f :: 'a :: \text{finite} \Rightarrow \text{complex}) = \text{sqrt } (\sum_{x \in UNIV}. (\text{norm } (f\ x))^2)$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-finite-L2-set*: $\text{ell2-norm } (x :: 'a :: \text{finite} \Rightarrow \text{complex}) = L2\text{-set } (\text{norm } o\ x) UNIV$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-square*: $\langle (\text{ell2-norm } x)^2 = (\sum_{\infty} i. (\text{cmod } (x\ i))^2) \rangle$
 $\langle \text{proof} \rangle$

lemma *ell2-ket*:
fixes a
defines $\langle f \equiv (\lambda i. \text{of-bool } (a = i)) \rangle$
shows *has-ell2-norm-ket*: $\langle \text{has-ell2-norm } f \rangle$
and *ell2-norm-ket*: $\langle \text{ell2-norm } f = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-geq0*: $\langle \text{ell2-norm } x \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-point-bound*:
assumes $\langle \text{has-ell2-norm } x \rangle$
shows $\langle \text{ell2-norm } x \geq \text{cmod } (x \ i) \rangle$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-0*:
assumes *has-ell2-norm* x
shows $\text{ell2-norm } x = 0 \longleftrightarrow x = (\lambda -. 0)$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-smult*:
assumes *has-ell2-norm* x
shows *has-ell2-norm* $(\lambda i. c * x \ i)$ **and** $\text{ell2-norm } (\lambda i. c * x \ i) = \text{cmod } c * \text{ell2-norm } x$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-triangle*:
assumes *has-ell2-norm* x **and** *has-ell2-norm* y
shows *has-ell2-norm* $(\lambda i. x \ i + y \ i)$ **and** $\text{ell2-norm } (\lambda i. x \ i + y \ i) \leq \text{ell2-norm } x + \text{ell2-norm } y$
 $\langle \text{proof} \rangle$

lemma *ell2-norm-uminus*:
assumes *has-ell2-norm* x
shows $\langle \text{has-ell2-norm } (\lambda i. - x \ i) \rangle$ **and** $\langle \text{ell2-norm } (\lambda i. - x \ i) = \text{ell2-norm } x \rangle$
 $\langle \text{proof} \rangle$

14.2 The type *ell2* of square-summable functions

typedef *'a ell2* = $\langle \{f :: 'a \Rightarrow \text{complex}. \text{has-ell2-norm } f\} \rangle$
 $\langle \text{proof} \rangle$

setup-lifting *type-definition-ell2*

instantiation *ell2* :: *(type)complex-vector* **begin**

lift-definition *zero-ell2* :: *'a ell2* **is** $\lambda -. 0$ $\langle \text{proof} \rangle$

lift-definition *uminus-ell2* :: *'a ell2* $\Rightarrow$ *'a ell2* **is** *uminus* $\langle \text{proof} \rangle$

lift-definition *plus-ell2* :: $\langle 'a \text{ ell2} \Rightarrow 'a \text{ ell2} \Rightarrow 'a \text{ ell2} \rangle$ **is** $\langle \lambda f \ g \ x. f \ x + g \ x \rangle$
 $\langle \text{proof} \rangle$

lift-definition *minus-ell2* :: $\langle 'a \text{ ell2} \Rightarrow 'a \text{ ell2} \Rightarrow 'a \text{ ell2} \rangle$ **is** $\lambda f \ g \ x. f \ x - g \ x$
 $\langle \text{proof} \rangle$

lift-definition *scaleR-ell2* :: *real* $\Rightarrow$ *'a ell2* $\Rightarrow$ *'a ell2* **is** $\lambda r \ f \ x. \text{complex-of-real } r * f \ x$
 $\langle \text{proof} \rangle$

lift-definition *scaleC-ell2* :: $\langle \text{complex} \Rightarrow 'a \text{ ell2} \Rightarrow 'a \text{ ell2} \rangle$ **is** $\langle \lambda c \ f \ x. c * f \ x \rangle$

$\langle \text{proof} \rangle$
instance
 $\langle \text{proof} \rangle$
end

instantiation $\text{ell2} :: (\text{type}) \text{ complex-normed-vector}$ **begin**
lift-definition $\text{norm-ell2} :: 'a \text{ ell2} \Rightarrow \text{real}$ **is** ell2-norm $\langle \text{proof} \rangle$
declare norm-ell2-def $[\text{code del}]$
definition $\text{dist } x \ y = \text{norm } (x - y)$ **for** $x \ y :: 'a \text{ ell2}$
definition $\text{sgn } x = x \ /_R \text{ norm } x$ **for** $x :: 'a \text{ ell2}$
definition $[\text{code del}]$: $\text{uniformity} = (\text{INF } e \in \{0 < ..\}. \text{ principal } \{(x :: 'a \text{ ell2}, y). \text{ norm } (x - y) < e\})$
definition $[\text{code del}]$: $\text{open } U = (\forall x \in U. \forall_F (x', y) \text{ in } \text{INF } e \in \{0 < ..\}. \text{ principal } \{(x, y). \text{ norm } (x - y) < e\}. x' = x \longrightarrow y \in U)$ **for** $U :: 'a \text{ ell2 set}$
instance
 $\langle \text{proof} \rangle$
end

lemma $\text{norm-point-bound-ell2}$: $\text{norm } (\text{Rep-ell2 } x \ i) \leq \text{norm } x$
 $\langle \text{proof} \rangle$

lemma $\text{ell2-norm-finite-support}$:
assumes $\langle \text{finite } S \rangle \ \langle \bigwedge i. i \notin S \implies \text{Rep-ell2 } x \ i = 0 \rangle$
shows $\langle \text{norm } x = \text{sqrt } ((\text{sum } (\lambda i. (\text{cmod } (\text{Rep-ell2 } x \ i))^2)) \ S) \rangle$
 $\langle \text{proof} \rangle$

instantiation $\text{ell2} :: (\text{type}) \text{ complex-inner}$ **begin**
lift-definition $\text{cinner-ell2} :: \langle 'a \text{ ell2} \Rightarrow 'a \text{ ell2} \Rightarrow \text{complex} \rangle$ **is**
 $\langle \lambda f \ g. \sum_{\infty} x. \text{cnj } (f \ x) * g \ x \rangle$ $\langle \text{proof} \rangle$
declare cinner-ell2-def $[\text{code del}]$

instance
 $\langle \text{proof} \rangle$
end

instance $\text{ell2} :: (\text{type}) \text{ hilbert-space}$
 $\langle \text{proof} \rangle$

lemma sum-ell2-transfer $[\text{transfer-rule}]$:
includes lifting-syntax
shows $\langle (((=) ==> \text{pcr-ell2 } (=)) ==> \text{rel-set } (=) ==> \text{pcr-ell2 } (=))$
 $\quad (\lambda f \ X \ x. \text{sum } (\lambda y. f \ y \ x) \ X) \text{ sum} \rangle$
 $\langle \text{proof} \rangle$

lemma clinear-Rep-ell2 $[\text{simp}]$: $\langle \text{clinear } (\lambda \psi. \text{Rep-ell2 } \psi \ i) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{Abs-ell2-inverse-finite}$ $[\text{simp}]$: $\langle \text{Rep-ell2 } (\text{Abs-ell2 } \psi) = \psi \rangle$ **for** $\psi :: \langle \cdot :: \text{finite} \rangle$

$\Rightarrow \text{complex}$
 $\langle \text{proof} \rangle$

14.3 Orthogonality

lemma *ell2-pointwise-ortho*:
assumes $\langle \bigwedge i. \text{Rep-ell2 } x \ i = 0 \ \vee \ \text{Rep-ell2 } y \ i = 0 \rangle$
shows $\langle \text{is-orthogonal } x \ y \rangle$
 $\langle \text{proof} \rangle$

14.4 Truncated vectors

lift-definition *trunc-ell2*:: $\langle 'a \ \text{set} \Rightarrow 'a \ \text{ell2} \Rightarrow 'a \ \text{ell2} \rangle$
is $\langle \lambda \ S \ x. (\lambda \ i. (\text{if } i \in S \text{ then } x \ i \text{ else } 0)) \rangle$
 $\langle \text{proof} \rangle$

lemma *trunc-ell2-empty[simp]*: $\langle \text{trunc-ell2 } \{\} \ x = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *trunc-ell2-UNIV[simp]*: $\langle \text{trunc-ell2 } \text{UNIV} \ \psi = \psi \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-id-minus-trunc-ell2*:
 $\langle (\text{norm } (x - \text{trunc-ell2 } S \ x))^2 = (\text{norm } x)^2 - (\text{norm } (\text{trunc-ell2 } S \ x))^2 \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-trunc-ell2-finite*:
 $\langle \text{finite } S \implies (\text{norm } (\text{trunc-ell2 } S \ x)) = \text{sqrt } ((\text{sum } (\lambda i. (\text{cmod } (\text{Rep-ell2 } x \ i))^2)) \ S) \rangle$
 $\langle \text{proof} \rangle$

lemma *trunc-ell2-lim-at-UNIV*:
 $\langle ((\lambda S. \text{trunc-ell2 } S \ \psi) \longrightarrow \psi) \ (\text{finite-subsets-at-top } \text{UNIV}) \rangle$
 $\langle \text{proof} \rangle$

lemma *trunc-ell2-lim-seq*: $\langle ((\lambda n. \text{trunc-ell2 } \{..<n\} \ \psi) \longrightarrow \psi) \rangle$
 $\langle \text{proof} \rangle$

lemma *trunc-ell2-norm-mono*: $\langle M \subseteq N \implies \text{norm } (\text{trunc-ell2 } M \ \psi) \leq \text{norm } (\text{trunc-ell2 } N \ \psi) \rangle$
 $\langle \text{proof} \rangle$

lemma *trunc-ell2-reduces-norm*: $\langle \text{norm } (\text{trunc-ell2 } M \ \psi) \leq \text{norm } \psi \rangle$
 $\langle \text{proof} \rangle$

lemma *trunc-ell2-twice[simp]*: $\langle \text{trunc-ell2 } M \ (\text{trunc-ell2 } N \ \psi) = \text{trunc-ell2 } (M \cap N) \ \psi \rangle$
 $\langle \text{proof} \rangle$

lemma *trunc-ell2-union*: $\langle \text{trunc-ell2 } (M \cup N) \psi = \text{trunc-ell2 } M \psi + \text{trunc-ell2 } N \psi - \text{trunc-ell2 } (M \cap N) \psi \rangle$

$\langle \text{proof} \rangle$

lemma *trunc-ell2-union-disjoint*: $\langle M \cap N = \{\} \implies \text{trunc-ell2 } (M \cup N) \psi = \text{trunc-ell2 } M \psi + \text{trunc-ell2 } N \psi \rangle$

$\langle \text{proof} \rangle$

lemma *trunc-ell2-union-Diff*: $\langle M \subseteq N \implies \text{trunc-ell2 } (N - M) \psi = \text{trunc-ell2 } N \psi - \text{trunc-ell2 } M \psi \rangle$

$\langle \text{proof} \rangle$

lemma *trunc-ell2-add*: $\langle \text{trunc-ell2 } M (\psi + \varphi) = \text{trunc-ell2 } M \psi + \text{trunc-ell2 } M \varphi \rangle$

$\langle \text{proof} \rangle$

lemma *trunc-ell2-scaleC*: $\langle \text{trunc-ell2 } M (c *_C \psi) = c *_C \text{trunc-ell2 } M \psi \rangle$

$\langle \text{proof} \rangle$

lemma *bounded-clinear-trunc-ell2*[*bounded-clinear*]: $\langle \text{bounded-clinear } (\text{trunc-ell2 } M) \rangle$

$\langle \text{proof} \rangle$

lemma *trunc-ell2-lim*: $\langle ((\lambda S. \text{trunc-ell2 } S \psi) \longrightarrow \text{trunc-ell2 } M \psi) \text{ (finite-subsets-at-top } M) \rangle$

$\langle \text{proof} \rangle$

lemma *trunc-ell2-lim-general*:

assumes *big*: $\langle \bigwedge G. \text{finite } G \implies G \subseteq M \implies (\forall_F H \text{ in } F. H \supseteq G) \rangle$

assumes *small*: $\langle \forall_F H \text{ in } F. H \subseteq M \rangle$

shows $\langle ((\lambda S. \text{trunc-ell2 } S \psi) \longrightarrow \text{trunc-ell2 } M \psi) F \rangle$

$\langle \text{proof} \rangle$

lemma *norm-ell2-bound-trunc*:

assumes $\langle \bigwedge M. \text{finite } M \implies \text{norm } (\text{trunc-ell2 } M \psi) \leq B \rangle$

shows $\langle \text{norm } \psi \leq B \rangle$

$\langle \text{proof} \rangle$

lemma *trunc-ell2-uminus*: $\langle \text{trunc-ell2 } (-M) \psi = \psi - \text{trunc-ell2 } M \psi \rangle$

$\langle \text{proof} \rangle$

14.5 Kets and bras

lift-definition *ket* :: $\langle 'a \Rightarrow 'a \text{ ell2} \rangle$ **is** $\langle \lambda x y. \text{of-bool } (x=y) \rangle$

$\langle \text{proof} \rangle$

abbreviation *bra* :: $\langle 'a \Rightarrow (-, \text{complex}) \text{ cblinfun} \rangle$ **where** *bra* *i* $\equiv$ *vector-to-cblinfun* (*ket i*)* **for** *i*

instance *ell2* :: (type) *not-singleton*

$\langle \text{proof} \rangle$

lemma *cinner-ket-left*: $\langle \text{ket } i \cdot_C \psi = \text{Rep-ell2 } \psi \ i \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-ket-right*: $\langle (\psi \cdot_C \text{ket } i) = \text{cnj } (\text{Rep-ell2 } \psi \ i) \rangle$
 $\langle \text{proof} \rangle$

lemma *bounded-clinear-Rep-ell2*[*simp*, *bounded-clinear*]: $\langle \text{bounded-clinear } (\lambda \psi. \text{Rep-ell2 } \psi \ x) \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-ket-eqI*:
 assumes $\langle \bigwedge i. \text{ket } i \cdot_C \psi = \text{ket } i \cdot_C \varphi \rangle$
 shows $\langle \psi = \varphi \rangle$
 $\langle \text{proof} \rangle$

lemma *norm-ket*[*simp*]: $\text{norm } (\text{ket } i) = 1$
 $\langle \text{proof} \rangle$

lemma *cinner-ket-same*[*simp*]:
 $\langle (\text{ket } i \cdot_C \text{ket } i) = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *orthogonal-ket*[*simp*]:
 $\langle \text{is-orthogonal } (\text{ket } i) \ (\text{ket } j) \longleftrightarrow i \neq j \rangle$
 $\langle \text{proof} \rangle$

lemma *cinner-ket*: $\langle (\text{ket } i \cdot_C \text{ket } j) = \text{of-bool } (i=j) \rangle$
 $\langle \text{proof} \rangle$

lemma *ket-injective*[*simp*]: $\langle \text{ket } i = \text{ket } j \longleftrightarrow i = j \rangle$
 $\langle \text{proof} \rangle$

lemma *inj-ket*[*simp*]: $\langle \text{inj-on } \text{ket } M \rangle$
 $\langle \text{proof} \rangle$

lemma *trunc-ell2-ket-cspan*:
 $\langle \text{trunc-ell2 } S \ x \in \text{cspan } (\text{range ket}) \rangle$ **if** $\langle \text{finite } S \rangle$
 $\langle \text{proof} \rangle$

lemma *closed-cspan-range-ket*[*simp*]:
 $\langle \text{closure } (\text{cspan } (\text{range ket})) = \text{UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *ccspan-range-ket*[*simp*]: $\text{ccspan } (\text{range ket}) = (\text{top}::('a \ \text{ell2} \ \text{ccsubspace}))$
 $\langle \text{proof} \rangle$

lemma *cspan-range-ket-finite*[simp]: *cspan* (*range ket* :: 'a::finite ell2 set) = UNIV
 ⟨proof⟩

instance *ell2* :: (finite) cfinite-dim
 ⟨proof⟩

instantiation *ell2* :: (enum) onb-enum **begin**

definition *canonical-basis-ell2* = map ket Enum.enum

definition ⟨*canonical-basis-length-ell2* (- :: 'a ell2 itself) = length (Enum.enum :: 'a list)⟩

instance
 ⟨proof⟩

end

lemma *canonical-basis-length-ell2*[code-unfold, simp]:
 length (*canonical-basis* :: 'a::enum ell2 list) = CARD('a)
 ⟨proof⟩

lemma *ket-canonical-basis*: ket *x* = *canonical-basis* ! *enum-idx x*
 ⟨proof⟩

lemma *clinear-equal-ket*:
 fixes *f g* :: ⟨'a::finite ell2 ⇒ -⟩
 assumes ⟨*clinear f*⟩
 assumes ⟨*clinear g*⟩
 assumes ⟨ $\bigwedge i. f \text{ (ket } i) = g \text{ (ket } i)$ ⟩
 shows ⟨*f* = *g*⟩
 ⟨proof⟩

lemma *equal-ket*:
 fixes *A B* :: ⟨('a ell2, 'b::complex-normed-vector) cblinfun⟩
 assumes ⟨ $\bigwedge x. A *_V \text{ ket } x = B *_V \text{ ket } x$ ⟩
 shows ⟨*A* = *B*⟩
 ⟨proof⟩

lemma *antilinear-equal-ket*:
 fixes *f g* :: ⟨'a::finite ell2 ⇒ -⟩
 assumes ⟨*antilinear f*⟩
 assumes ⟨*antilinear g*⟩
 assumes ⟨ $\bigwedge i. f \text{ (ket } i) = g \text{ (ket } i)$ ⟩
 shows ⟨*f* = *g*⟩
 ⟨proof⟩

lemma *cinner-ket-adjointI*:
 fixes *F*::'a ell2 ⇒_{CL} - and *G*::'b ell2 ⇒_{CL} -
 assumes $\bigwedge i j. (F *_V \text{ ket } i) \cdot_C \text{ ket } j = \text{ket } i \cdot_C (G *_V \text{ ket } j)$
 shows *F* = *G**
 ⟨proof⟩

lemma *ket-nonzero*[simp]: $\text{ket } i \neq 0$
 ⟨proof⟩

lemma *cindependent-ket*[simp]:
 $\text{cindependent } (\text{range } (\text{ket}::'a \Rightarrow -))$
 ⟨proof⟩

lemma *cdim-UNIV-ell2*[simp]: $\langle \text{cdim } (\text{UNIV}::'a::\text{finite ell2 set}) = \text{CARD}('a) \rangle$
 ⟨proof⟩

lemma *is-ortho-set-ket*[simp]: $\langle \text{is-ortho-set } (\text{range ket}) \rangle$
 ⟨proof⟩

lemma *bounded-clinear-equal-ket*:
 fixes $f g :: 'a \text{ ell2} \Rightarrow -$
 assumes $\langle \text{bounded-clinear } f \rangle$
 assumes $\langle \text{bounded-clinear } g \rangle$
 assumes $\langle \bigwedge i. f (\text{ket } i) = g (\text{ket } i) \rangle$
 shows $\langle f = g \rangle$
 ⟨proof⟩

lemma *bounded-antilinear-equal-ket*:
 fixes $f g :: 'a \text{ ell2} \Rightarrow -$
 assumes $\langle \text{bounded-antilinear } f \rangle$
 assumes $\langle \text{bounded-antilinear } g \rangle$
 assumes $\langle \bigwedge i. f (\text{ket } i) = g (\text{ket } i) \rangle$
 shows $\langle f = g \rangle$
 ⟨proof⟩

lemma *is-onb-ket*[simp]: $\langle \text{is-onb } (\text{range ket}) \rangle$
 ⟨proof⟩

lemma *ell2-sum-ket*: $\langle \psi = (\sum_{i \in \text{UNIV}. \text{Rep-ell2 } \psi \ i *_{\text{C}} \text{ket } i}) \rangle$ **for** $\psi :: \langle -::\text{finite ell2} \rangle$
 ⟨proof⟩

lemma *trunc-ell2-singleton*: $\langle \text{trunc-ell2 } \{x\} \psi = \text{Rep-ell2 } \psi \ x *_{\text{C}} \text{ket } x \rangle$
 ⟨proof⟩

lemma *trunc-ell2-insert*: $\langle \text{trunc-ell2 } (\text{insert } x \ M) \ \varphi = \text{Rep-ell2 } \varphi \ x *_{\text{C}} \text{ket } x + \text{trunc-ell2 } M \ \varphi \rangle$
if $\langle x \notin M \rangle$
 ⟨proof⟩

lemma *trunc-ell2-finite-sum*: $\langle \text{trunc-ell2 } M \ \psi = (\sum_{i \in M}. \text{Rep-ell2 } \psi \ i *_{\text{C}} \text{ket } i) \rangle$
if $\langle \text{finite } M \rangle$
 ⟨proof⟩

lemma *is-orthogonal-trunc-ell2*: $\langle \text{is-orthogonal } (\text{trunc-ell2 } M \ \psi) \ (\text{trunc-ell2 } N \ \varphi) \rangle$

if $\langle M \cap N = \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma *separating-set-ket*: $\langle \text{separating-set bounded-clinear } (\text{range ket}) \rangle$
 $\langle \text{proof} \rangle$

14.6 Butterflies

lemma *cspan-butterfly-ket*: $\langle \text{cspan } \{ \text{butterfly } (\text{ket } i) (\text{ket } j) \mid (i::'b::\text{finite}) (j::'a::\text{finite}). \text{True} \} = \text{UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *cindependent-butterfly-ket*: $\langle \text{cindependent } \{ \text{butterfly } (\text{ket } i) (\text{ket } j) \mid (i::'b) (j::'a). \text{True} \} \rangle$
 $\langle \text{proof} \rangle$

lemma *clinear-eq-butterfly-ketI*:
fixes $F \ G :: \langle ('a::\text{finite ell2} \Rightarrow_{CL} 'b::\text{finite ell2}) \Rightarrow 'c::\text{complex-vector} \rangle$
assumes *clinear F and clinear G*
assumes $\bigwedge i \ j. F (\text{butterfly } (\text{ket } i) (\text{ket } j)) = G (\text{butterfly } (\text{ket } i) (\text{ket } j))$
shows $F = G$
 $\langle \text{proof} \rangle$

lemma *sum-butterfly-ket[simp]*: $\langle (\sum (i::'a::\text{finite}) \in \text{UNIV}. \text{butterfly } (\text{ket } i) (\text{ket } i)) = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *ell2-decompose-has-sum*: $\langle ((\lambda x. \text{Rep-ell2 } \varphi \ x \ *_C \text{ket } x) \text{ has-sum } \varphi) \text{ UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *ell2-decompose-infsum*: $\langle \varphi = (\sum_{\infty} x. \text{Rep-ell2 } \varphi \ x \ *_C \text{ket } x) \rangle$
 $\langle \text{proof} \rangle$

lemma *ell2-decompose-summable*: $\langle (\lambda x. \text{Rep-ell2 } \varphi \ x \ *_C \text{ket } x) \text{ summable-on UNIV} \rangle$
 $\langle \text{proof} \rangle$

lemma *Rep-ell2-cblinfun-apply-sum*: $\langle \text{Rep-ell2 } (A \ *_V \varphi) \ y = (\sum_{\infty} x. \text{Rep-ell2 } \varphi \ x \ *_V \text{ket } x) \ y \rangle$
 $\langle \text{proof} \rangle$

14.7 One-dimensional spaces

instantiation *ell2* :: (*CARD-1*) **one begin**
lift-definition *one-ell2* :: 'a *ell2* **is** $\lambda-. 1$ $\langle \text{proof} \rangle$
instance $\langle \text{proof} \rangle$
end

lemma *ket-CARD-1-is-1*: $\langle \text{ket } x = 1 \rangle$ **for** $x :: \langle 'a::\text{CARD-1} \rangle$
 $\langle \text{proof} \rangle$

instantiation *ell2* :: (*CARD-1*) *times* **begin**
lift-definition *times-ell2* :: '*a ell2* $\Rightarrow$ '*a ell2* $\Rightarrow$ '*a ell2* **is** $\lambda a\ b\ x. a\ x * b\ x$
 $\langle proof \rangle$
instance $\langle proof \rangle$
end

instantiation *ell2* :: (*CARD-1*) *divide* **begin**
lift-definition *divide-ell2* :: '*a ell2* $\Rightarrow$ '*a ell2* $\Rightarrow$ '*a ell2* **is** $\lambda a\ b\ x. a\ x / b\ x$
 $\langle proof \rangle$
instance $\langle proof \rangle$
end

instantiation *ell2* :: (*CARD-1*) *inverse* **begin**
lift-definition *inverse-ell2* :: '*a ell2* $\Rightarrow$ '*a ell2* **is** $\lambda a\ x. inverse\ (a\ x)$
 $\langle proof \rangle$
instance $\langle proof \rangle$
end

instance *ell2* :: (*{enum, CARD-1}*) *one-dim*

Note: *enum* is not needed logically, but without it this instantiation clashes with *instantiation ell2 :: (enum) onb-enum*

$\langle proof \rangle$

14.8 Explicit bounded operators

definition *explicit-cblinfun* :: $\langle ('a \Rightarrow 'b \Rightarrow complex) \Rightarrow ('b\ ell2, 'a\ ell2)\ cblinfun \rangle$
where
 $\langle explicit-cblinfun\ M = cblinfun-extension\ (range\ ket)\ (\lambda a. Abs-ell2\ (\lambda j. M\ j\ (inv\ ket\ a))) \rangle$

definition *explicit-cblinfun-exists* :: $\langle ('a \Rightarrow 'b \Rightarrow complex) \Rightarrow bool \rangle$ **where**
 $\langle explicit-cblinfun-exists\ M \longleftrightarrow$
 $(\forall a. has-ell2-norm\ (\lambda j. M\ j\ a)) \wedge$
 $cblinfun-extension-exists\ (range\ ket)\ (\lambda a. Abs-ell2\ (\lambda j. M\ j\ (inv\ ket\ a))) \rangle$

lemma *explicit-cblinfun-exists-bounded*:

assumes $\langle \bigwedge S\ T\ \psi. finite\ S \implies finite\ T \implies (\bigwedge a. a \notin T \implies \psi\ a = 0) \implies$
 $(\sum b \in S. (cmod\ (\sum a \in T. \psi\ a *_{\mathcal{C}}\ M\ b\ a))^2) \leq B * (\sum a \in T. (cmod\ (\psi$
 $a))^2 \rangle$
shows $\langle explicit-cblinfun-exists\ M \rangle$
 $\langle proof \rangle$

lemma *explicit-cblinfun-exists-finite-dim[simp]*: $\langle explicit-cblinfun-exists\ m \rangle$ **for** *m*
 $:: \neg finite \Rightarrow \neg finite \Rightarrow -$
 $\langle proof \rangle$

lemma *explicit-cblinfun-ket*: $\langle explicit-cblinfun\ M *_{\mathcal{V}}\ ket\ a = Abs-ell2\ (\lambda b. M\ b\ a) \rangle$
if $\langle explicit-cblinfun-exists\ M \rangle$

⟨proof⟩

lemma *Rep-ell2-explicit-cblinfun-ket[simp]*: $\langle \text{Rep-ell2 } (\text{explicit-cblinfun } M *_{\mathcal{V}} \text{ket } a) \text{ } b = M \text{ } b \text{ } a \rangle \text{ if } \langle \text{explicit-cblinfun-exists } M \rangle$
 ⟨proof⟩

lemma *bounded-extension-counterexample-1*: $\langle \exists f. \forall x. f (\text{ket } x) = \text{ket } 0 \rangle$
 — First part of counterexample showing that not every linear function can be extended to a bounded operator.
 ⟨proof⟩

lemma *bounded-extension-counterexample-2*:
 — Second part of counterexample showing that not every linear function can be extended to a bounded operator.
assumes $\langle \forall x::'a::\text{infinite}. f (\text{ket } x) = \text{ket } 0 \rangle$
shows $\langle \neg \text{cblinfun-extension-exists } (\text{range } \text{ket}) \text{ } f \rangle$
 ⟨proof⟩

14.9 Diagonal operators

definition *diagonal-operator* **where** $\langle \text{diagonal-operator } f =$
 (if bdd-above (range ($\lambda x. \text{cmod } (f \text{ } x)$))) then explicit-cblinfun ($\lambda x \text{ } y. \text{of-bool } (x=y)$
 $* f \text{ } x$) else 0) $\rangle$

lemma *diagonal-operator-exists*:
assumes $\langle \text{bdd-above } (\text{range } (\lambda x. \text{cmod } (f \text{ } x))) \rangle$
shows $\langle \text{explicit-cblinfun-exists } (\lambda x \text{ } y. \text{of-bool } (x = y) * f \text{ } x) \rangle$
 ⟨proof⟩

lemma *diagonal-operator-ket*:
assumes $\langle \text{bdd-above } (\text{range } (\lambda x. \text{cmod } (f \text{ } x))) \rangle$
shows $\langle \text{diagonal-operator } f (\text{ket } x) = f \text{ } x *_C \text{ket } x \rangle$
 ⟨proof⟩

lemma *diagonal-operator-invalid*:
assumes $\langle \neg \text{bdd-above } (\text{range } (\lambda x. \text{cmod } (f \text{ } x))) \rangle$
shows $\langle \text{diagonal-operator } f = 0 \rangle$
 ⟨proof⟩

lemma *diagonal-operator-adj*: $\langle \text{diagonal-operator } f^* = \text{diagonal-operator } (\lambda x. \text{cnj } (f \text{ } x)) \rangle$
 ⟨proof⟩

lemma *diagonal-operator-comp*:
assumes $\langle \text{bdd-above } (\text{range } (\lambda x. \text{cmod } (f \text{ } x))) \rangle$
assumes $\langle \text{bdd-above } (\text{range } (\lambda x. \text{cmod } (g \text{ } x))) \rangle$

shows $\langle \text{diagonal-operator } f \circ_{CL} \text{diagonal-operator } g = \text{diagonal-operator } (\lambda x. (f x * g x)) \rangle$
 $\langle \text{proof} \rangle$

14.10 Classical operators

We call an operator mapping $\text{ket } x$ to $\text{ket } (\pi x)$ or 0 "classical". (The meaning is inspired by the fact that in quantum mechanics, such operators usually correspond to operations with classical interpretation (such as Pauli-X, CNOT, measurement in the computational basis, etc.))

definition $\text{classical-operator} :: ('a \Rightarrow 'b \text{ option}) \Rightarrow 'a \text{ ell2} \Rightarrow_{CL} 'b \text{ ell2}$ **where**
 $\text{classical-operator } \pi =$
 $(\text{let } f = (\lambda t. (\text{case } \pi (\text{inv } (\text{ket}::'a \Rightarrow -)) t$
 $\quad \text{of None} \Rightarrow (0::'b \text{ ell2})$
 $\quad | \text{Some } i \Rightarrow \text{ket } i))$
in
 $\text{cblinfun-extension } (\text{range } (\text{ket}::'a \Rightarrow -)) f)$

definition $\text{classical-operator-exists } \pi \longleftrightarrow$
 $\text{cblinfun-extension-exists } (\text{range } \text{ket})$
 $(\lambda t. \text{case } \pi (\text{inv } \text{ket } t) \text{ of None} \Rightarrow 0 \mid \text{Some } i \Rightarrow \text{ket } i)$

lemma $\text{classical-operator-existsI}$:
assumes $\bigwedge x. B *_V (\text{ket } x) = (\text{case } \pi x \text{ of Some } i \Rightarrow \text{ket } i \mid \text{None} \Rightarrow 0)$
shows $\text{classical-operator-exists } \pi$
 $\langle \text{proof} \rangle$

lemma
assumes $\text{inj-map } \pi$
shows $\text{classical-operator-exists-inj}$: $\text{classical-operator-exists } \pi$
and $\text{classical-operator-norm-inj}$: $\langle \text{norm } (\text{classical-operator } \pi) \leq 1 \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{classical-operator-exists-finite[simp]}$: $\text{classical-operator-exists } (\pi :: \text{finite} \Rightarrow -)$
 $\langle \text{proof} \rangle$

lemma $\text{classical-operator-ket}$:
assumes $\text{classical-operator-exists } \pi$
shows $(\text{classical-operator } \pi) *_V (\text{ket } x) = (\text{case } \pi x \text{ of Some } i \Rightarrow \text{ket } i \mid \text{None} \Rightarrow 0)$
 $\langle \text{proof} \rangle$

lemma $\text{classical-operator-ket-finite}$:
 $(\text{classical-operator } \pi) *_V (\text{ket } (x::'a::\text{finite})) = (\text{case } \pi x \text{ of Some } i \Rightarrow \text{ket } i \mid \text{None} \Rightarrow 0)$
 $\langle \text{proof} \rangle$

```

lemma classical-operator-adjoint[simp]:
  fixes  $\pi :: 'a \Rightarrow 'b \text{ option}$ 
  assumes  $a1: \text{inj-map } \pi$ 
  shows  $(\text{classical-operator } \pi)^* = \text{classical-operator } (\text{inv-map } \pi)$ 
   $\langle \text{proof} \rangle$ 

lemma
  fixes  $\pi :: 'b \Rightarrow 'c \text{ option}$  and  $\varrho :: 'a \Rightarrow 'b \text{ option}$ 
  assumes classical-operator-exists  $\pi$ 
  assumes classical-operator-exists  $\varrho$ 
  shows classical-operator-exists-comp[simp]: classical-operator-exists  $(\pi \circ_m \varrho)$ 
    and classical-operator-mult[simp]: classical-operator  $\pi \circ_{CL}$  classical-operator  $\varrho$ 
   $= \text{classical-operator } (\pi \circ_m \varrho)$ 
   $\langle \text{proof} \rangle$ 

lemma classical-operator-Some[simp]: classical-operator  $(\text{Some} :: 'a \Rightarrow -) = \text{id-cblinfun}$ 
   $\langle \text{proof} \rangle$ 

lemma isometry-classical-operator[simp]:
  fixes  $\pi :: 'a \Rightarrow 'b$ 
  assumes  $a1: \text{inj } \pi$ 
  shows isometry  $(\text{classical-operator } (\text{Some } o \pi))$ 
   $\langle \text{proof} \rangle$ 

lemma unitary-classical-operator[simp]:
  fixes  $\pi :: 'a \Rightarrow 'b$ 
  assumes  $a1: \text{bij } \pi$ 
  shows unitary  $(\text{classical-operator } (\text{Some } o \pi))$ 
   $\langle \text{proof} \rangle$ 

unbundle no lattice-syntax and no cblinfun-syntax

end

```

15 *Extra-Jordan-Normal-Form* – Additional results for `Jordan_Normal_Form`

```

theory Extra-Jordan-Normal-Form
  imports
    Jordan-Normal-Form.Matrix Jordan-Normal-Form.Schur-Decomposition
begin

```

We define bundles to activate/deactivate the notation from `Jordan_Normal_Form`. Reactivate the notation locally via "**includes** *jnf-syntax*" in a lemma statement. (Or sandwich a declaration using that notation between "**unbundle** *jnf-syntax* ... **unbundle** *no jnf-syntax*.)

```

open-bundle jnf-syntax
begin

```

```

notation transpose-mat ( $\langle (-^T) \rangle$  [1000])
notation cscalar-prod (infix  $\langle \cdot c \rangle$  70)
notation vec-index (infixl  $\langle \$ \rangle$  100)
notation smult-vec (infixl  $\langle \cdot_v \rangle$  70)
notation scalar-prod (infix  $\langle \cdot \rangle$  70)
notation index-mat (infixl  $\langle \$\$ \rangle$  100)
notation smult-mat (infixl  $\langle \cdot_m \rangle$  70)
notation mult-mat-vec (infixl  $\langle *_v \rangle$  70)
notation pow-mat (infixr  $\langle \hat{\cdot}_m \rangle$  75)
notation append-vec (infixr  $\langle @_v \rangle$  65)
notation append-rows (infixr  $\langle @_r \rangle$  65)
end

```

```

lemma mat-entry-explicit:
  fixes  $M :: 'a::field\ mat$ 
  assumes  $M \in carrier\_mat\ m\ n$  and  $i < m$  and  $j < n$ 
  shows  $vec\_index\ (M *_v\ unit\_vec\ n\ j)\ i = M\ \$\$ (i,j)$ 
   $\langle proof \rangle$ 

```

```

lemma mat-adjoint-def':  $mat\_adjoint\ M = transpose\_mat\ (map\_mat\ conjugate\ M)$ 
   $\langle proof \rangle$ 

```

```

lemma mat-adjoint-swap:
  fixes  $M :: complex\ mat$ 
  assumes  $M \in carrier\_mat\ nB\ nA$  and  $iA < dim\_row\ M$  and  $iB < dim\_col\ M$ 
  shows  $(mat\_adjoint\ M)\ \$\$ (iB,iA) = cnj\ (M\ \$\$ (iA,iB))$ 
   $\langle proof \rangle$ 

```

```

lemma cscalar-prod-adjoint:
  fixes  $M :: complex\ mat$ 
  assumes  $M \in carrier\_mat\ nB\ nA$ 
  and  $dim\_vec\ v = nA$ 
  and  $dim\_vec\ u = nB$ 
  shows  $v \cdot c\ ((mat\_adjoint\ M) *_v\ u) = (M *_v\ v) \cdot c\ u$ 
   $\langle proof \rangle$ 

```

```

lemma scaleC-minus1-left-vec:  $-1 \cdot_v v = -\ v$  for  $v :: ring-1\ vec$ 
   $\langle proof \rangle$ 

```

```

lemma square-nneg-complex:
  fixes  $x :: complex$ 
  assumes  $x \in \mathbb{R}$  shows  $x^2 \geq 0$ 
   $\langle proof \rangle$ 

```

```

definition vec-is-zero  $n\ v = (\forall i < n. v\ \$\ i = 0)$ 

```

```

lemma vec-is-zero:  $dim\_vec\ v = n \implies vec\_is\_zero\ n\ v \longleftrightarrow v = 0_v\ n$ 
   $\langle proof \rangle$ 

```



```

fun gram-schmidt-sub0
  where gram-schmidt-sub0 n us [] = us
  | gram-schmidt-sub0 n us (w # ws) =
    (let w' = adjuster n w us + w in
     if vec-is-zero n w' then gram-schmidt-sub0 n us ws
     else gram-schmidt-sub0 n (w' # us) ws)

```

```

lemma (in cof-vec-space) adjuster-already-in-span:
  assumes w ∈ carrier-vec n
  assumes us-carrier: set us ⊆ carrier-vec n
  assumes corthogonal us
  assumes w ∈ span (set us)
  shows adjuster n w us + w = 0_v n
⟨proof⟩

```

```

lemma (in cof-vec-space) gram-schmidt-sub0-result:
  assumes gram-schmidt-sub0 n us ws = us'
  and set ws ⊆ carrier-vec n
  and set us ⊆ carrier-vec n
  and distinct us
  and ~ lin-dep (set us)
  and corthogonal us
  shows set us' ⊆ carrier-vec n ∧
    distinct us' ∧
    corthogonal us' ∧
    span (set (us @ ws)) = span (set us')
⟨proof⟩

```

This is a variant of *gram-schmidt* that does not require the input vectors *ws* to be distinct or linearly independent. (In comparison to *gram-schmidt*, our version also returns the result in reversed order.)

definition *gram-schmidt0* n ws = *gram-schmidt-sub0* n [] ws

```

lemma (in cof-vec-space) gram-schmidt0-result:
  fixes ws
  defines us' ≡ gram-schmidt0 n ws
  assumes ws: set ws ⊆ carrier-vec n
  shows set us' ⊆ carrier-vec n (is ?thesis1)
  and distinct us' (is ?thesis2)
  and corthogonal us' (is ?thesis3)
  and span (set ws) = span (set us') (is ?thesis4)
⟨proof⟩

```

locale *complex-vec-space* = *cof-vec-space* n TYPE(*complex*) **for** n :: nat

```

lemma gram-schmidt0-corthogonal:
  assumes a1: corthogonal R

```

and $a2: \bigwedge x. x \in \text{set } R \implies \text{dim-vec } x = d$
shows $\text{gram-schmidt0 } d \ R = \text{rev } R$
 $\langle \text{proof} \rangle$

lemma *adjuster-carrier'*:
assumes $w: (w :: 'a::\text{conjugatable-field vec}) : \text{carrier-vec } n$
and $us: \text{set } (us :: 'a \text{ vec list}) \subseteq \text{carrier-vec } n$
shows $\text{adjuster } n \ w \ us \in \text{carrier-vec } n$
 $\langle \text{proof} \rangle$

lemma *eq-mat-on-vecI*:
fixes $M \ N :: \langle 'a::\text{field mat} \rangle$
assumes $\text{eq}: \langle \bigwedge v. v \in \text{carrier-vec } nA \implies M *_v v = N *_v v \rangle$
assumes $[\text{simp}]: \langle M \in \text{carrier-mat } nB \ nA \rangle \langle N \in \text{carrier-mat } nB \ nA \rangle$
shows $\langle M = N \rangle$
 $\langle \text{proof} \rangle$

lemma *list-of-vec-plus*:
fixes $v1 \ v2 :: \langle \text{complex vec} \rangle$
assumes $\langle \text{dim-vec } v1 = \text{dim-vec } v2 \rangle$
shows $\langle \text{list-of-vec } (v1 + v2) = \text{map2 } (+) (\text{list-of-vec } v1) (\text{list-of-vec } v2) \rangle$
 $\langle \text{proof} \rangle$

lemma *list-of-vec-mult*:
fixes $v :: \langle \text{complex vec} \rangle$
shows $\langle \text{list-of-vec } (c \cdot_v v) = \text{map } ((*) \ c) (\text{list-of-vec } v) \rangle$
 $\langle \text{proof} \rangle$

lemma *map-map-vec-cols*: $\langle \text{map } (\text{map-vec } f) (\text{cols } m) = \text{cols } (\text{map-mat } f \ m) \rangle$
 $\langle \text{proof} \rangle$

lemma *map-vec-conjugate*: $\langle \text{map-vec conjugate } v = \text{conjugate } v \rangle$
 $\langle \text{proof} \rangle$

unbundle *no jnf-syntax*

end

16 Cblinfun-Matrix – Matrix representation of bounded operators

theory *Cblinfun-Matrix*

imports

Complex-L2

Jordan-Normal-Form.Gram-Schmidt

HOL–Analysis.Starlike

Complex-Bounded-Operators.Extra-Jordan-Normal-Form

begin

hide-const (**open**) *Order.bottom Order.top*
hide-type (**open**) *Finite-Cartesian-Product.vec*
hide-const (**open**) *Finite-Cartesian-Product.mat*
hide-fact (**open**) *Finite-Cartesian-Product.mat-def*
hide-const (**open**) *Finite-Cartesian-Product.vec*
hide-fact (**open**) *Finite-Cartesian-Product.vec-def*
hide-const (**open**) *Finite-Cartesian-Product.row*
hide-fact (**open**) *Finite-Cartesian-Product.row-def*
no-notation *Finite-Cartesian-Product.vec-nth* (**infixl** $\langle \$ \rangle$ 90)

unbundle *jnf-syntax*
unbundle *cblinfun-syntax*

16.1 Isomorphism between vectors

We define the canonical isomorphism between vectors in some complex vector space $'a$ and the complex n -dimensional vectors (where n is the dimension of $'a$). This is possible if $'a$, $'b$ are of class *basis-enum* since that class fixes a finite canonical basis. Vector are represented using the *complex vec* type from *Jordan_Normal_Form*. (The isomorphism will be called *vec-of-onb-basis* below.)

definition *vec-of-basis-enum* :: $\langle 'a::\text{basis-enum} \Rightarrow \text{complex vec} \rangle$ **where**
 — Maps v to a $'a$ *vec* represented in basis *canonical-basis*
 $\langle \text{vec-of-basis-enum } v = \text{vec-of-list } (\text{map } (\text{crepresentation } (\text{set canonical-basis}) v) \text{ canonical-basis}) \rangle$

lemma *dim-vec-of-basis-enum*[*simp*]:
 $\langle \text{dim-vec } (\text{vec-of-basis-enum } (v::'a)) = \text{length } (\text{canonical-basis}::'a::\text{basis-enum list}) \rangle$
 $\langle \text{proof} \rangle$

definition *basis-enum-of-vec* :: $\langle \text{complex vec} \Rightarrow 'a::\text{basis-enum} \rangle$ **where**
 $\langle \text{basis-enum-of-vec } v =$
 (if $\text{dim-vec } v = \text{length } (\text{canonical-basis}::'a \text{ list})$
 then $\text{sum-list } (\text{map2 } (*_C) (\text{list-of-vec } v) (\text{canonical-basis}::'a \text{ list}))$
 else 0) $\rangle$

lemma *vec-of-basis-enum-inverse*[*simp*]:
fixes $\psi :: 'a::\text{basis-enum}$
shows $\text{basis-enum-of-vec } (\text{vec-of-basis-enum } \psi) = \psi$
 $\langle \text{proof} \rangle$

lemma *basis-enum-of-vec-inverse*[*simp*]:
fixes $v :: \text{complex vec}$
defines $n \equiv \text{length } (\text{canonical-basis}::'a::\text{basis-enum list})$
assumes $f1: \text{dim-vec } v = n$
shows $\text{vec-of-basis-enum } ((\text{basis-enum-of-vec } v)::'a) = v$

$\langle \text{proof} \rangle$

lemma *basis-enum-eq-vec-of-basis-enumI*:
fixes $a\ b :: \text{basis-enum}$
assumes $\text{vec-of-basis-enum } a = \text{vec-of-basis-enum } b$
shows $a = b$
 $\langle \text{proof} \rangle$

lemma *vec-of-basis-enum-carrier-vec[simp]*: $\langle \text{vec-of-basis-enum } v \in \text{carrier-vec } (\text{canonical-basis-length } \text{TYPE}('a)) \rangle$ **for** $v :: \langle 'a :: \text{basis-enum} \rangle$
 $\langle \text{proof} \rangle$

lemma *vec-of-basis-enum-inj*: $\text{inj } \text{vec-of-basis-enum}$
 $\langle \text{proof} \rangle$

lemma *basis-enum-of-vec-inj*: $\text{inj-on } (\text{basis-enum-of-vec} :: \text{complex } \text{vec} \Rightarrow 'a)$
 $(\text{carrier-vec } (\text{length } (\text{canonical-basis} :: 'a :: \{\text{basis-enum}, \text{complex-normed-vector}\}$
 $\text{list})))$
 $\langle \text{proof} \rangle$

16.2 Operations on vectors

lemma *basis-enum-of-vec-add*:
assumes $[\text{simp}]$: $\langle \text{dim-vec } v1 = \text{length } (\text{canonical-basis} :: 'a :: \text{basis-enum } \text{list}) \rangle$
 $\langle \text{dim-vec } v2 = \text{length } (\text{canonical-basis} :: 'a \text{ list}) \rangle$
shows $\langle ((\text{basis-enum-of-vec } (v1 + v2)) :: 'a) = \text{basis-enum-of-vec } v1 + \text{basis-enum-of-vec } v2 \rangle$
 $\langle \text{proof} \rangle$

lemma *basis-enum-of-vec-mult*:
assumes $[\text{simp}]$: $\langle \text{dim-vec } v = \text{length } (\text{canonical-basis} :: 'a :: \text{basis-enum } \text{list}) \rangle$
shows $\langle ((\text{basis-enum-of-vec } (c \cdot_v v)) :: 'a) = c *_C \text{basis-enum-of-vec } v \rangle$
 $\langle \text{proof} \rangle$

lemma *vec-of-basis-enum-add*:
 $\langle \text{vec-of-basis-enum } (a + b) = \text{vec-of-basis-enum } a + \text{vec-of-basis-enum } b \rangle$
 $\langle \text{proof} \rangle$

lemma *vec-of-basis-enum-scaleC*:
 $\text{vec-of-basis-enum } (c *_C b) = c \cdot_v (\text{vec-of-basis-enum } b)$
 $\langle \text{proof} \rangle$

lemma *vec-of-basis-enum-scaleR*:
 $\text{vec-of-basis-enum } (r *_R b) = \text{complex-of-real } r \cdot_v (\text{vec-of-basis-enum } b)$
 $\langle \text{proof} \rangle$

lemma *vec-of-basis-enum-uminus*:
 $\text{vec-of-basis-enum } (- b2) = - \text{vec-of-basis-enum } b2$
 $\langle \text{proof} \rangle$

lemma *vec-of-basis-enum-minus*:

vec-of-basis-enum (*b1* - *b2*) = *vec-of-basis-enum* *b1* - *vec-of-basis-enum* *b2*
 ⟨*proof*⟩

lemma *cinner-basis-enum-of-vec*:

defines *n* ≡ *length* (*canonical-basis* :: 'a::onb-enum list)
assumes [*simp*]: *dim-vec* *x* = *n* *dim-vec* *y* = *n*
shows (*basis-enum-of-vec* *x* :: 'a) •_C *basis-enum-of-vec* *y* = *y* •_C *x*
 ⟨*proof*⟩

lemma *cscalar-prod-vec-of-basis-enum*: *cscalar-prod* (*vec-of-basis-enum* *φ*) (*vec-of-basis-enum* *ψ*) = *cinner* *ψ* *φ*

for *ψ* :: 'a::onb-enum
 ⟨*proof*⟩

definition ⟨*norm-vec* *ψ* = *sqrt* ($\sum i \in \{0 \dots \dim\text{-vec } \psi\}$). let *z* = *vec-index* *ψ* *i* in (*Re* *z*)² + (*Im* *z*)²)

lemma *norm-vec-of-basis-enum*: ⟨*norm* *ψ* = *norm-vec* (*vec-of-basis-enum* *ψ*) **for** *ψ* :: 'a::onb-enum
 ⟨*proof*⟩

lemma *basis-enum-of-vec-unit-vec*:

defines *basis* ≡ (*canonical-basis* :: 'a::basis-enum list)
and *n* ≡ *length* (*canonical-basis* :: 'a list)
assumes *a3*: *i* < *n*
shows *basis-enum-of-vec* (*unit-vec* *n* *i*) = *basis*!*i*
 ⟨*proof*⟩

lemma *vec-of-basis-enum-ket*:

vec-of-basis-enum (*ket* *i*) = *unit-vec* (*CARD*('a)) (*enum-idx* *i*)
for *i*::'a::enum
 ⟨*proof*⟩

lemma *vec-of-basis-enum-zero*:

defines *nA* ≡ *length* (*canonical-basis* :: 'a::basis-enum list)
shows *vec-of-basis-enum* (0::'a) = 0_v *nA*
 ⟨*proof*⟩

lemma (in *complex-vec-space*) *vec-of-basis-enum-cspan*:

fixes *X* :: 'a::basis-enum set
assumes *length* (*canonical-basis* :: 'a list) = *n*
shows *vec-of-basis-enum* ' *cspan* *X* = *span* (*vec-of-basis-enum* ' *X*)
 ⟨*proof*⟩

lemma (in *complex-vec-space*) *basis-enum-of-vec-span*:

assumes *length* (*canonical-basis* :: 'a list) = *n*
assumes *Y* ⊆ *carrier-vec* *n*

shows $\text{basis-enum-of-vec } \langle \text{local.span } Y = \text{cspan } (\text{basis-enum-of-vec } \langle Y :: 'a::\text{basis-enum set} \rangle) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{vec-of-basis-enum-canonical-basis}$:
assumes $i < \text{length } (\text{canonical-basis} :: 'a \text{ list})$
shows $\text{vec-of-basis-enum } (\text{canonical-basis}!i :: 'a)$
 $= \text{unit-vec } (\text{length } (\text{canonical-basis} :: 'a::\text{basis-enum list})) i$
 $\langle \text{proof} \rangle$

lemma $\text{vec-of-basis-enum-times}$:
fixes $\psi \ \varphi :: 'a::\text{one-dim}$
shows $\text{vec-of-basis-enum } (\psi * \varphi)$
 $= \text{vec-of-list } [\text{vec-index } (\text{vec-of-basis-enum } \psi) \ 0 * \text{vec-index } (\text{vec-of-basis-enum } \varphi) \ 0]$
 $\langle \text{proof} \rangle$

lemma $\text{vec-of-basis-enum-to-inverse}$:
fixes $\psi :: 'a::\text{one-dim}$
shows $\text{vec-of-basis-enum } (\text{inverse } \psi) = \text{vec-of-list } [\text{inverse } (\text{vec-index } (\text{vec-of-basis-enum } \psi) \ 0)]$
 $\langle \text{proof} \rangle$

lemma $\text{vec-of-basis-enum-divide}$:
fixes $\psi \ \varphi :: 'a::\text{one-dim}$
shows $\text{vec-of-basis-enum } (\psi / \varphi)$
 $= \text{vec-of-list } [\text{vec-index } (\text{vec-of-basis-enum } \psi) \ 0 / \text{vec-index } (\text{vec-of-basis-enum } \varphi) \ 0]$
 $\langle \text{proof} \rangle$

lemma $\text{vec-of-basis-enum-1}$: $\text{vec-of-basis-enum } (1 :: 'a::\text{one-dim}) = \text{vec-of-list } [1]$
 $\langle \text{proof} \rangle$

lemma $\text{vec-of-basis-enum-ell2-component}$:
fixes $\psi :: \langle 'a::\text{enum ell2} \rangle$
assumes $[\text{simp}]$: $\langle i < \text{CARD}('a) \rangle$
shows $\langle \text{vec-of-basis-enum } \psi \ \$ \ i = \text{Rep-ell2 } \psi \ (\text{Enum.enum } ! \ i) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{corthogonal-vec-of-basis-enum}$:
fixes $S :: 'a::\text{onb-enum list}$
shows $\text{corthogonal } (\text{map } \text{vec-of-basis-enum } S) \longleftrightarrow \text{is-ortho-set } (\text{set } S) \wedge \text{distinct } S$
 $\langle \text{proof} \rangle$

16.3 Isomorphism between bounded linear functions and matrices

We define the canonical isomorphism between $'a \Rightarrow_{CL} 'b$ and the complex $n * m$ -matrices (where n, m are the dimensions of $'a, 'b$, respectively). This is possible if $'a, 'b$ are of class *basis-enum* since that class fixes a finite canonical basis. Matrices are represented using the *complex mat* type from *Jordan_Normal_Form*. (The isomorphism will be called *mat-of-cblinfun* below.)

definition *mat-of-cblinfun* :: $\langle 'a :: \{ \text{basis-enum}, \text{complex-normed-vector} \} \Rightarrow_{CL} 'b :: \{ \text{basis-enum}, \text{complex-normed-vector} \} \Rightarrow \text{complex mat} \rangle$ **where**
 $\langle \text{mat-of-cblinfun } f =$
 $\text{mat } (\text{length } (\text{canonical-basis} :: 'b \text{ list})) (\text{length } (\text{canonical-basis} :: 'a \text{ list})) ($
 $\lambda (i, j). \text{crepresentation } (\text{set } (\text{canonical-basis} :: 'b \text{ list})) (f *_{\mathcal{V}} ((\text{canonical-basis} :: 'a$
 $\text{list})!j)) ((\text{canonical-basis} :: 'b \text{ list})!i)) \rangle$
for f

lift-definition *cblinfun-of-mat* :: $\langle \text{complex mat} \Rightarrow 'a :: \{ \text{basis-enum}, \text{complex-normed-vector} \} \Rightarrow_{CL} 'b :: \{ \text{basis-enum}, \text{complex-normed-vector} \} \rangle$ **is**
 $\langle \lambda M. \text{if } M \in \text{carrier-mat } (\text{length } (\text{canonical-basis} :: 'b \text{ list})) (\text{length } (\text{canonical-basis}$
 $:: 'a \text{ list}))$
 $\text{then } \lambda v. \text{basis-enum-of-vec } (M *_{\mathcal{V}} \text{vec-of-basis-enum } v)$
 $\text{else } (\lambda v. 0) \rangle$
 $\langle \text{proof} \rangle$

lemma *cblinfun-of-mat-invalid*:

assumes $\langle M \notin \text{carrier-mat } (\text{canonical-basis-length TYPE('b :: \{ basis-enum, complex-normed-vector \})})$
 $(\text{canonical-basis-length TYPE('a :: \{ basis-enum, complex-normed-vector \})}) \rangle$
shows $\langle (\text{cblinfun-of-mat } M :: 'a \Rightarrow_{CL} 'b) = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *dim-row-mat-of-cblinfun[simp]*: $\langle \text{dim-row } (\text{mat-of-cblinfun } (a :: 'a :: \{ \text{basis-enum}, \text{complex-normed-vector} \} \Rightarrow_{CL} 'b :: \{ \text{basis-enum}, \text{complex-normed-vector} \})) = \text{canonical-basis-length TYPE('b)} \rangle$
 $\langle \text{proof} \rangle$

lemma *dim-col-mat-of-cblinfun[simp]*: $\langle \text{dim-col } (\text{mat-of-cblinfun } (a :: 'a :: \{ \text{basis-enum}, \text{complex-normed-vector} \} \Rightarrow_{CL} 'b :: \{ \text{basis-enum}, \text{complex-normed-vector} \})) = \text{canonical-basis-length TYPE('a)} \rangle$
 $\langle \text{proof} \rangle$

lemma *mat-of-cblinfun-ell2-carrier[simp]*: $\langle \text{mat-of-cblinfun } (a :: 'a :: \{ \text{basis-enum}, \text{complex-normed-vector} \} \Rightarrow_{CL} 'b :: \{ \text{basis-enum}, \text{complex-normed-vector} \}) \in \text{carrier-mat } (\text{canonical-basis-length TYPE('b)}) (\text{canonical-basis-length TYPE('a)}) \rangle$
 $\langle \text{proof} \rangle$

lemma *basis-enum-of-vec-cblinfun-apply*:

fixes $M :: \text{complex mat}$
defines $nA \equiv \text{length } (\text{canonical-basis} :: 'a :: \{ \text{basis-enum}, \text{complex-normed-vector} \} \text{ list})$

and $nB \equiv \text{length } (\text{canonical-basis} :: 'b::\{\text{basis-enum}, \text{complex-normed-vector}\} \text{ list})$
assumes $M \in \text{carrier-mat } nB \ nA$ **and** $\text{dim-vec } x = nA$
shows $\text{basis-enum-of-vec } (M *_V x) = (\text{cblinfun-of-mat } M :: 'a \Rightarrow_{CL} 'b) *_V \text{basis-enum-of-vec } x$
 $\langle \text{proof} \rangle$

lemma *mat-of-cblinfun-cblinfun-apply*:
 $\langle \text{vec-of-basis-enum } (F *_V u) = \text{mat-of-cblinfun } F *_V \text{vec-of-basis-enum } u \rangle$
for $F::'a::\{\text{basis-enum}, \text{complex-normed-vector}\} \Rightarrow_{CL} 'b::\{\text{basis-enum}, \text{complex-normed-vector}\}$
and $u::'a$
 $\langle \text{proof} \rangle$

lemma *mat-of-cblinfun-inverse*:
 $\text{cblinfun-of-mat } (\text{mat-of-cblinfun } B) = B$
for $B::'a::\{\text{basis-enum}, \text{complex-normed-vector}\} \Rightarrow_{CL} 'b::\{\text{basis-enum}, \text{complex-normed-vector}\}$
 $\langle \text{proof} \rangle$

lemma *mat-of-cblinfun-inj*: $\text{inj mat-of-cblinfun}$
 $\langle \text{proof} \rangle$

lemma *cblinfun-of-mat-inverse*:
fixes $M::\text{complex mat}$
defines $nA \equiv \text{length } (\text{canonical-basis} :: 'a::\{\text{basis-enum}, \text{complex-normed-vector}\} \text{ list})$
and $nB \equiv \text{length } (\text{canonical-basis} :: 'b::\{\text{basis-enum}, \text{complex-normed-vector}\} \text{ list})$
assumes $M \in \text{carrier-mat } nB \ nA$
shows $\text{mat-of-cblinfun } (\text{cblinfun-of-mat } M :: 'a \Rightarrow_{CL} 'b) = M$
 $\langle \text{proof} \rangle$

lemma *cblinfun-of-mat-inj*: $\text{inj-on } (\text{cblinfun-of-mat}::\text{complex mat} \Rightarrow 'a \Rightarrow_{CL} 'b)$
 $(\text{carrier-mat } (\text{length } (\text{canonical-basis} :: 'b::\{\text{basis-enum}, \text{complex-normed-vector}\} \text{ list})))$
 $(\text{length } (\text{canonical-basis} :: 'a::\{\text{basis-enum}, \text{complex-normed-vector}\} \text{ list})))$
 $\langle \text{proof} \rangle$

lemma *cblinfun-eq-mat-of-cblinfunI*:
assumes $\text{mat-of-cblinfun } a = \text{mat-of-cblinfun } b$
shows $a = b$
 $\langle \text{proof} \rangle$

16.4 Operations on matrices

lemma *cblinfun-of-mat-plus*:
defines $nA \equiv \text{length } (\text{canonical-basis} :: 'a::\{\text{basis-enum}, \text{complex-normed-vector}\} \text{ list})$
and $nB \equiv \text{length } (\text{canonical-basis} :: 'b::\{\text{basis-enum}, \text{complex-normed-vector}\} \text{ list})$

list)
assumes [simp,intro]: $M \in \text{carrier-mat } nB \ nA$ **and** [simp,intro]: $N \in \text{carrier-mat } nB \ nA$
shows (cblinfun-of-mat ($M + N$) :: ' $a \Rightarrow_{CL} 'b$) = ((cblinfun-of-mat M + cblinfun-of-mat N))
 <proof>

lemma mat-of-cblinfun-zero:

mat-of-cblinfun ($0 :: ('a::\{\text{basis-enum, complex-normed-vector}\} \Rightarrow_{CL} 'b::\{\text{basis-enum, complex-normed-vector}\})$
 = 0_m (length (canonical-basis :: ' b list)) (length (canonical-basis :: ' a list))
 <proof>

lemma mat-of-cblinfun-plus:

mat-of-cblinfun ($F + G$) = mat-of-cblinfun F + mat-of-cblinfun G
for $F \ G::'a::\{\text{basis-enum, complex-normed-vector}\} \Rightarrow_{CL} 'b::\{\text{basis-enum, complex-normed-vector}\}$
 <proof>

lemma mat-of-cblinfun-id:

mat-of-cblinfun (id-cblinfun :: (' $a::\{\text{basis-enum, complex-normed-vector}\} \Rightarrow_{CL} 'a$))
 = 1_m (length (canonical-basis :: ' a list))
 <proof>

lemma mat-of-cblinfun-1:

mat-of-cblinfun ($1 :: ('a::\text{one-dim} \Rightarrow_{CL} 'b::\text{one-dim})$) = $1_m \ 1$
 <proof>

lemma mat-of-cblinfun-uminus:

mat-of-cblinfun ($- M$) = $-$ mat-of-cblinfun M
for $M::'a::\{\text{basis-enum, complex-normed-vector}\} \Rightarrow_{CL} 'b::\{\text{basis-enum, complex-normed-vector}\}$
 <proof>

lemma mat-of-cblinfun-minus:

mat-of-cblinfun ($M - N$) = mat-of-cblinfun M - mat-of-cblinfun N
for $M::'a::\{\text{basis-enum, complex-normed-vector}\} \Rightarrow_{CL} 'b::\{\text{basis-enum, complex-normed-vector}\}$
and $N::'a \Rightarrow_{CL} 'b$
 <proof>

lemma cblinfun-of-mat-uminus:

defines $nA \equiv \text{length (canonical-basis :: 'a::\{\text{basis-enum, complex-normed-vector}\} list)}$
and $nB \equiv \text{length (canonical-basis :: 'b::\{\text{basis-enum, complex-normed-vector}\} list)}$
assumes $M \in \text{carrier-mat } nB \ nA$
shows (cblinfun-of-mat ($-M$) :: ' $a \Rightarrow_{CL} 'b$) = $-$ cblinfun-of-mat M
 <proof>

lemma cblinfun-of-mat-minus:

fixes $M::\text{complex mat}$
defines $nA \equiv \text{length (canonical-basis :: 'a::\{\text{basis-enum, complex-normed-vector}\} list)}$

$list)$
and $nB \equiv \text{length } (canonical\text{-}basis :: 'b::\{basis\text{-}enum, complex\text{-}normed\text{-}vector\})$
 $list)$
assumes $M \in carrier\text{-}mat\ nB\ nA$ **and** $N \in carrier\text{-}mat\ nB\ nA$
shows $(cblinfun\text{-}of\text{-}mat\ (M - N) :: 'a \Rightarrow_{CL} 'b) = cblinfun\text{-}of\text{-}mat\ M - cblinfun\text{-}of\text{-}mat\ N$
 $\langle proof \rangle$

lemma *cblinfun-of-mat-times*:
fixes $M\ N :: complex\ mat$
defines $nA \equiv \text{length } (canonical\text{-}basis :: 'a::\{basis\text{-}enum, complex\text{-}normed\text{-}vector\})$
 $list)$
and $nB \equiv \text{length } (canonical\text{-}basis :: 'b::\{basis\text{-}enum, complex\text{-}normed\text{-}vector\})$
 $list)$
and $nC \equiv \text{length } (canonical\text{-}basis :: 'c::\{basis\text{-}enum, complex\text{-}normed\text{-}vector\})$
 $list)$
assumes $a1: M \in carrier\text{-}mat\ nC\ nB$ **and** $a2: N \in carrier\text{-}mat\ nB\ nA$
shows $cblinfun\text{-}of\text{-}mat\ (M * N) = ((cblinfun\text{-}of\text{-}mat\ M) :: 'b \Rightarrow_{CL} 'c) \circ_{CL} ((cblinfun\text{-}of\text{-}mat\ N) :: 'a \Rightarrow_{CL} 'b)$
 $\langle proof \rangle$

lemma *cblinfun-of-mat-adjoint*:
defines $nA \equiv \text{length } (canonical\text{-}basis :: 'a::onb\text{-}enum\ list)$
and $nB \equiv \text{length } (canonical\text{-}basis :: 'b::onb\text{-}enum\ list)$
fixes $M :: complex\ mat$
assumes $M \in carrier\text{-}mat\ nB\ nA$
shows $((cblinfun\text{-}of\text{-}mat\ (mat\text{-}adjoint\ M)) :: 'b \Rightarrow_{CL} 'a) = (cblinfun\text{-}of\text{-}mat\ M)^*$
 $\langle proof \rangle$

lemma *mat-of-cblinfun-compose*:
 $mat\text{-}of\text{-}cblinfun\ (F \circ_{CL} G) = mat\text{-}of\text{-}cblinfun\ F * mat\text{-}of\text{-}cblinfun\ G$
for $F :: 'b::\{basis\text{-}enum, complex\text{-}normed\text{-}vector\} \Rightarrow_{CL} 'c::\{basis\text{-}enum, complex\text{-}normed\text{-}vector\}$
and $G :: 'a::\{basis\text{-}enum, complex\text{-}normed\text{-}vector\} \Rightarrow_{CL} 'b$
 $\langle proof \rangle$

lemma *mat-of-cblinfun-scaleC*:
 $mat\text{-}of\text{-}cblinfun\ ((a::complex) *_C F) = a \cdot_m (mat\text{-}of\text{-}cblinfun\ F)$
for $F :: 'a::\{basis\text{-}enum, complex\text{-}normed\text{-}vector\} \Rightarrow_{CL} 'b::\{basis\text{-}enum, complex\text{-}normed\text{-}vector\}$
 $\langle proof \rangle$

lemma *mat-of-cblinfun-scaleR*:
 $mat\text{-}of\text{-}cblinfun\ ((a::real) *_R F) = (complex\text{-}of\text{-}real\ a) \cdot_m (mat\text{-}of\text{-}cblinfun\ F)$
 $\langle proof \rangle$

lemma *mat-of-cblinfun-adj*:
 $mat\text{-}of\text{-}cblinfun\ (F^*) = mat\text{-}adjoint\ (mat\text{-}of\text{-}cblinfun\ F)$
for $F :: 'a::onb\text{-}enum \Rightarrow_{CL} 'b::onb\text{-}enum$
 $\langle proof \rangle$

lemma *mat-of-cblinfun-vector-to-cblinfun*:

```
mat-of-cblinfun (vector-to-cblinfun  $\psi$ )
  = mat-of-cols (length (canonical-basis :: 'a list)) [vec-of-basis-enum  $\psi$ ]
for  $\psi :: 'a :: \{basis-enum, complex-normed-vector\}$ 
<proof>
```

lemma *mat-of-cblinfun-proj*:

```
fixes  $a :: 'a :: onb-enum$ 
defines  $d \equiv length (canonical-basis :: 'a list)$ 
and  $norm2 \equiv (vec-of-basis-enum a) \cdot c (vec-of-basis-enum a)$ 
shows  $mat-of-cblinfun (proj a) =$ 
   $1 / norm2 \cdot_m (mat-of-cols d [vec-of-basis-enum a]$ 
     $* mat-of-rows d [conjugate (vec-of-basis-enum a)])$  (is  $\leftarrow = ?rhs$ )
<proof>
```

lemma *mat-of-cblinfun-ell2-component*:

```
fixes  $a :: \langle 'a :: enum\ ell2 \Rightarrow_{CL} 'b :: enum\ ell2 \rangle$ 
assumes [simp]:  $\langle i < CARD('b) \rangle \langle j < CARD('a) \rangle$ 
shows  $\langle mat-of-cblinfun a \$\$ (i,j) = Rep-ell2 (a *_V ket (Enum.enum ! j))$ 
   $(Enum.enum ! i) \rangle$ 
<proof>
```

lemma *cblinfun-of-mat-mat*:

```
shows  $\langle cblinfun-of-mat (mat (CARD('b)) (CARD('a)) f) = explicit-cblinfun$ 
   $(\lambda(r :: 'b :: enum) (c :: 'a :: enum). f (enum-idx r, enum-idx c)) \rangle$ 
<proof>
```

lemma *mat-of-cblinfun-explicit-cblinfun*:

```
fixes  $f :: \langle 'a :: enum \Rightarrow 'b :: enum \Rightarrow complex \rangle$ 
shows  $\langle mat-of-cblinfun (explicit-cblinfun f) = mat (CARD('a)) (CARD('b))$ 
   $(\lambda(r,c). f (Enum.enum!r) (Enum.enum!c)) \rangle$ 
<proof>
```

lemma *mat-of-cblinfun-classical-operator*:

```
fixes  $f :: 'a :: enum \Rightarrow 'b :: enum\ option$ 
shows  $mat-of-cblinfun (classical-operator f) = mat (CARD('b)) (CARD('a))$ 
   $(\lambda(r,c). \text{if } f (Enum.enum!c) = \text{Some } (Enum.enum!r) \text{ then } 1 \text{ else } 0)$ 
<proof>
```

lemma *mat-of-cblinfun-sandwich*:

```
fixes  $a :: (- :: onb-enum, - :: onb-enum) cblinfun$ 
shows  $\langle mat-of-cblinfun (sandwich a *_V b) = (let a' = mat-of-cblinfun a \text{ in } a' *$ 
   $mat-of-cblinfun b * mat-adjoint a') \rangle$ 
<proof>
```

lemma *mat-of-cblinfun-one-dim-iso*:

```
 $\langle mat-of-cblinfun (one-dim-iso f :: 'a :: one-dim \Rightarrow_{CL} 'b :: one-dim) = mat-of-rows-list$ 
```

1 [[one-dim-iso f]]
 <proof>

lemma *mat-of-cblinfun-times*:
 fixes $F\ G :: \langle 'a::\text{one-dim} \Rightarrow_{CL} 'b::\text{one-dim} \rangle$
 shows $\langle \text{mat-of-cblinfun } (F * G) = \text{mat-of-rows-list } 1\ [[(\text{one-dim-iso } F) * (\text{one-dim-iso } G)]] \rangle$
 <proof>

16.5 Operations on subspaces

lemma *ccspan-gram-schmidt0-invariant*:
 defines $\text{basis-enum} \equiv (\text{basis-enum-of-vec} :: - \Rightarrow 'a::\{\text{basis-enum}, \text{complex-normed-vector}\})$
 defines $n \equiv \text{length } (\text{canonical-basis} :: 'a\ \text{list})$
 assumes $\text{set } ws \subseteq \text{carrier-vec } n$
 shows $\text{ccspan } (\text{set } (\text{map } \text{basis-enum } (\text{gram-schmidt0 } n\ ws))) = \text{ccspan } (\text{set } (\text{map } \text{basis-enum } ws))$
 <proof>

definition *is-subspace-of-vec-list* $n\ vs\ ws =$
 (let $ws' = \text{gram-schmidt0 } n\ ws$ in
 $\forall v \in \text{set } vs. \text{adjuster } n\ v\ ws' = -\ v$)

lemma *ccspan-leq-using-vec*:
 fixes $A\ B :: \langle 'a::\{\text{basis-enum}, \text{complex-normed-vector}\} \text{ list} \rangle$
 shows $\langle (\text{ccspan } (\text{set } A) \leq \text{ccspan } (\text{set } B)) \longleftrightarrow$
 $\text{is-subspace-of-vec-list } (\text{length } (\text{canonical-basis} :: 'a\ \text{list}))$
 $(\text{map } \text{vec-of-basis-enum } A) (\text{map } \text{vec-of-basis-enum } B) \rangle$
 <proof>

lemma *cblinfun-image-ccspan-using-vec*:
 $A *_S \text{ccspan } (\text{set } S) = \text{ccspan } (\text{basis-enum-of-vec } ' \text{set } (\text{map } ((*_v) (\text{mat-of-cblinfun } A)) (\text{map } \text{vec-of-basis-enum } S)))$
 <proof>

mk-projector-orthog $d\ L$ takes a list L of d -dimensional vectors and returns the projector onto the span of L . (Assuming that all vectors in L are orthogonal and nonzero.)

fun *mk-projector-orthog* $:: \text{nat} \Rightarrow \text{complex vec list} \Rightarrow \text{complex mat}$ **where**
 $\text{mk-projector-orthog } d\ [] = \text{zero-mat } d\ d$
 $| \text{mk-projector-orthog } d\ [v] = (\text{let } \text{norm2} = \text{cscalar-prod } v\ v \text{ in}$
 $\text{smult-mat } (1/\text{norm2}) (\text{mat-of-cols } d\ [v] * \text{mat-of-rows } d$
 $[\text{conjugate } v]))$
 $| \text{mk-projector-orthog } d\ (v\#\text{vs}) = (\text{let } \text{norm2} = \text{cscalar-prod } v\ v \text{ in}$
 $\text{smult-mat } (1/\text{norm2}) (\text{mat-of-cols } d\ [v] * \text{mat-of-rows}$
 $d\ [\text{conjugate } v])$
 $+ \text{mk-projector-orthog } d\ \text{vs})$

```

lemma mk-projector-orthog-correct:
  fixes  $S :: 'a::\text{onb-enum list}$ 
  defines  $d \equiv \text{length } (\text{canonical-basis } :: 'a \text{ list})$ 
  assumes ortho: is-ortho-set (set S) and distinct: distinct S
  shows mk-projector-orthog  $d$  (map vec-of-basis-enum S)
    = mat-of-cblinfun (Proj (ccspan (set S)))
  <proof>

definition <mk-projector d vs = mk-projector-orthog d (gram-schmidt0 d vs)>

lemma mat-of-cblinfun-Proj-ccspan:
  fixes  $S :: \langle 'a::\text{onb-enum list} \rangle$ 
  shows <mat-of-cblinfun (Proj (ccspan (set S))) = mk-projector (length (canonical-basis
  :: 'a list)) (map vec-of-basis-enum S)>
  <proof>

unbundle no jnf-syntax and no cblinfun-syntax

end

```

17 Cblinfun-Code – Support for code generation

This theory provides support for code generation involving on complex vector spaces and bounded operators (e.g., types *cblinfun* and *ell2*). To fully support code generation, in addition to importing this theory, one need to activate support for code generation (import theory *Jordan-Normal-Form.Matrix-Impl*) and for real and complex numbers (import theory *Real-Impl.Real-Impl* for support of reals of the form $a + b * \text{sqrt } c$ or *Algebraic-Numbers.Real-Factorization* (much slower) for support of algebraic reals; support of complex numbers comes "for free").

The builtin support for real and complex numbers (in *Complex-Main*) is not sufficient because it does not support the computation of square-roots which are used in the setup below.

It is also recommended to import *HOL-Library.Code-Target-Numeral* for faster support of nats and integers.

```

theory Cblinfun-Code
  imports
    Cblinfun-Matrix Containers.Set-Impl Jordan-Normal-Form.Matrix-Kernel
  begin

  no-notation Lattice.meet (infixl  $\langle \sqcap_1 \rangle$  70)
  no-notation Lattice.join (infixl  $\langle \sqcup_1 \rangle$  65)
  hide-const (open) Coset.kernel
  hide-const (open) Matrix-Kernel.kernel
  hide-const (open) Order.bottom Order.top

```

```

unbundle lattice-syntax
unbundle jnf-syntax
unbundle cblinfun-syntax

```

17.1 Code equations for cblinfun operators

In this subsection, we define the code for all operations involving only operators (no combinations of operators/vectors/subspaces)

The following lemma registers cblinfun as an abstract datatype with constructor *cblinfun-of-mat*. That means that in generated code, all cblinfun operators will be represented as *cblinfun-of-mat X* where X is a matrix. In code equations for operations involving operators (e.g., +), we can then write the equation directly in terms of matrices by writing, e.g., *mat-of-cblinfun (A + B)* in the lhs, and in the rhs we define the matrix that corresponds to the sum of A,B. In the rhs, we can access the matrices corresponding to A,B by writing *mat-of-cblinfun B*. (See, e.g., lemma *mat-of-cblinfun-plus*.) See [2] for more information on [code abstype].

```
declare mat-of-cblinfun-inverse [code abstype]
```

```
declare mat-of-cblinfun-plus[code]
— Code equation for addition of cblinfun's
```

```
declare mat-of-cblinfun-id[code]
— Code equation for computing the identity operator
```

```
declare mat-of-cblinfun-1 [code]
— Code equation for computing the one-dimensional identity
```

```
declare mat-of-cblinfun-zero[code]
— Code equation for computing the zero operator
```

```
declare mat-of-cblinfun-uminus[code]
— Code equation for computing the unary minus on cblinfun's
```

```
declare mat-of-cblinfun-minus[code]
— Code equation for computing the difference of cblinfun's
```

```
declare mat-of-cblinfun-classical-operator[code]
— Code equation for computing the "classical operator"
```

```
declare mat-of-cblinfun-explicit-cblinfun[code]
— Code equation for computing the explicit-cblinfun
```

```
declare mat-of-cblinfun-compose[code]
```

— Code equation for computing the composition/product of cblinfun's

declare *mat-of-cblinfun-scaleC*[*code*]

— Code equation for multiplication with complex scalar

declare *mat-of-cblinfun-scaleR*[*code*]

— Code equation for multiplication with real scalar

declare *mat-of-cblinfun-adj*[*code*]

— Code equation for computing the adj

This instantiation defines a code equation for equality tests for cblinfun.

instantiation *cblinfun* :: (*onb-enum*, *onb-enum*) *equal* **begin**

definition [*code*]: *equal-cblinfun* *M N* $\longleftrightarrow$ *mat-of-cblinfun* *M* = *mat-of-cblinfun* *N*

for *M N* :: '*a* $\Rightarrow_{CL}$ 'b

instance

<proof>

end

17.2 Vectors

In this section, we define code for operations on vectors. As with operators above, we do this by using an isomorphism between finite vectors (i.e., types *T* of sort *complex-vector*) and the type *complex vec* from *Jordan_Normal_Form*. We have developed such an isomorphism in theory *Cblinfun-Matrix* for any type *T* of sort *onb-enum* (i.e., any type with a finite canonical orthonormal basis) as was done above for bounded operators. Unfortunately, we cannot declare code equations for a type class, code equations must be related to a specific type constructor. So we give code definition only for vectors of type '*a ell2* (where '*a* must be of sort *enum* to make make sure that '*a ell2* is finite dimensional).

The isomorphism between '*a ell2* is given by the constants *ell2-of-vec* and *vec-of-ell2* which are copies of the more general *basis-enum-of-vec* and *vec-of-basis-enum* but with a more restricted type to be usable in our code equations.

definition *ell2-of-vec* :: *complex vec* $\Rightarrow$ '*a::enum ell2* **where** *ell2-of-vec* = *basis-enum-of-vec*

definition *vec-of-ell2* :: '*a::enum ell2* $\Rightarrow$ *complex vec* **where** *vec-of-ell2* = *vec-of-basis-enum*

The following theorem registers the isomorphism *ell2-of-vec/vec-of-ell2* for code generation. From now on, code for operations on - *ell2* can be expressed by declarations such as *vec-of-ell2* (*f a b*) = *g* (*vec-of-ell2 a*) (*vec-of-ell2 b*) if the operation *f* on - *ell2* corresponds to the operation *g* on *complex vec*.

lemma *vec-of-ell2-inverse* [*code abstype*]:

ell2-of-vec (*vec-of-ell2 B*) = *B*

<proof>

This instantiation defines a code equation for equality tests for `ell2`.

instantiation `ell2` :: (*enum*) *equal* **begin**

definition [*code*]: *equal-ell2* *M N* $\longleftrightarrow$ *vec-of-ell2* *M* = *vec-of-ell2* *N*

for *M N* :: '*a*::*enum ell2*

instance

<proof>

end

lemma *vec-of-ell2-carrier-vec*[*simp*]: *<vec-of-ell2 v* $\in$ *carrier-vec* *CARD('a)* **for** *v*

:: '*a*::*enum ell2*

<proof>

lemma *vec-of-ell2-zero*[*code*]:

— Code equation for computing the zero vector

vec-of-ell2 (*0*::'*a*::*enum ell2*) = *zero-vec* (*CARD('a)*)

<proof>

lemma *vec-of-ell2-ket*[*code*]:

— Code equation for computing a standard basis vector

vec-of-ell2 (*ket i*) = *unit-vec* (*CARD('a)*) (*enum-idx i*)

for *i*::'*a*::*enum*

<proof>

lemma *vec-of-ell2-scaleC*[*code*]:

— Code equation for multiplying a vector with a complex scalar

vec-of-ell2 (*scaleC a* *ψ*) = *smult-vec* *a* (*vec-of-ell2* *ψ*)

for *ψ* :: '*a*::*enum ell2*

<proof>

lemma *vec-of-ell2-scaleR*[*code*]:

— Code equation for multiplying a vector with a real scalar

vec-of-ell2 (*scaleR a* *ψ*) = *smult-vec* (*complex-of-real a*) (*vec-of-ell2* *ψ*)

for *ψ* :: '*a*::*enum ell2*

<proof>

lemma *ell2-of-vec-plus*[*code*]:

— Code equation for adding vectors

vec-of-ell2 (*x* + *y*) = (*vec-of-ell2* *x*) + (*vec-of-ell2* *y*) **for** *x y* :: '*a*::*enum ell2*

<proof>

lemma *ell2-of-vec-minus*[*code*]:

— Code equation for subtracting vectors

vec-of-ell2 (*x* - *y*) = (*vec-of-ell2* *x*) - (*vec-of-ell2* *y*) **for** *x y* :: '*a*::*enum ell2*

<proof>

lemma *ell2-of-vec-uminus*[*code*]:

— Code equation for negating a vector

vec-of-ell2 (- *y*) = - (*vec-of-ell2* *y*) **for** *y* :: '*a*::*enum ell2*

<proof>

lemma *cinner-ell2-code* [code]: *cinner* ψ φ = *cscalar-prod* (*vec-of-ell2* φ) (*vec-of-ell2* ψ)

— Code equation for the inner product of vectors
 ⟨proof⟩

lemma *norm-ell2-code* [code]:

— Code equation for the norm of a vector
norm ψ = *norm-vec* (*vec-of-ell2* ψ)
 ⟨proof⟩

lemma *times-ell2-code*[code]:

— Code equation for the product in the algebra of one-dimensional vectors
fixes ψ φ :: 'a::{CARD-1,enum} *ell2*
shows *vec-of-ell2* ($\psi * \varphi$)
 = *vec-of-list* [*vec-index* (*vec-of-ell2* ψ) 0 * *vec-index* (*vec-of-ell2* φ) 0]
 ⟨proof⟩

lemma *divide-ell2-code*[code]:

— Code equation for the product in the algebra of one-dimensional vectors
fixes ψ φ :: 'a::{CARD-1,enum} *ell2*
shows *vec-of-ell2* (ψ / φ)
 = *vec-of-list* [*vec-index* (*vec-of-ell2* ψ) 0 / *vec-index* (*vec-of-ell2* φ) 0]
 ⟨proof⟩

lemma *inverse-ell2-code*[code]:

— Code equation for the product in the algebra of one-dimensional vectors
fixes ψ :: 'a::{CARD-1,enum} *ell2*
shows *vec-of-ell2* (*inverse* ψ)
 = *vec-of-list* [*inverse* (*vec-index* (*vec-of-ell2* ψ) 0)]
 ⟨proof⟩

lemma *one-ell2-code*[code]:

— Code equation for the unit in the algebra of one-dimensional vectors
vec-of-ell2 (1 :: 'a::{CARD-1,enum} *ell2*) = *vec-of-list* [1]
 ⟨proof⟩

17.3 Vector/Matrix

We proceed to give code equations for operations involving both operators (*cblinfun*) and vectors. As explained above, we have to restrict the equations to vectors of type 'a *ell2* even though the theory is available for any type of class *onb-enum*. As a consequence, we run into an additional technicality now. For example, to define a code equation for applying an operator to a vector, we might try to give the following lemma:

lemma *cblinfun-apply-ell2*[code]: *vec-of-ell2* ($M *_V x$) = (*mult-mat-vec* (*mat-of-cblinfun* M) (*vec-of-ell2* x)) **by** (*simp add: mat-of-cblinfun-cblinfun-apply vec-of-ell2-def*)

Unfortunately, this does not work, Isabelle produces the warning "Projection as head in equation", most likely due to the fact that the type of $(*_V)$ in the equation is less general than the type of $(*_V)$ (it is restricted to $ell2$). We overcome this problem by defining a constant *cblinfun-apply-ell2* which is equal to $(*_V)$ but has a more restricted type. We then instruct the code generation to replace occurrences of $(*_V)$ by *cblinfun-apply-ell2* (where possible), and we add code generation for *cblinfun-apply-ell2* instead of $(*_V)$.

definition *cblinfun-apply-ell2* :: '*a ell2* $\Rightarrow_{CL}$ '*b ell2* $\Rightarrow$ '*a ell2* $\Rightarrow$ '*b ell2*
where [*code del*, *code-abbrev*]: *cblinfun-apply-ell2* = $(*_V)$
— *code-abbrev* instructs the code generation to replace the rhs $(*_V)$ by the lhs *cblinfun-apply-ell2* before starting the actual code generation.

lemma *cblinfun-apply-ell2*[*code*]:
— Code equation for *cblinfun-apply-ell2*, i.e., for applying an operator to an *ell2* vector
 $vec\text{-of-ell2 } (cblinfun\text{-apply-ell2 } M \ x) = (mult\text{-mat-vec } (mat\text{-of-cblinfun } M) \ (vec\text{-of-ell2 } x))$
 $\langle proof \rangle$

For the constant *vector-to-cblinfun* (canonical isomorphism from vectors to operators), we have the same problem and define a constant *vector-to-cblinfun-code* with more restricted type

definition *vector-to-cblinfun-code* :: '*a ell2* $\Rightarrow$ '*b::one-dim* $\Rightarrow_{CL}$ '*a ell2* **where**
[*code del*, *code-abbrev*]: *vector-to-cblinfun-code* = *vector-to-cblinfun*
— *code-abbrev* instructs the code generation to replace the rhs *vector-to-cblinfun* by the lhs *vector-to-cblinfun-code* before starting the actual code generation.

lemma *vector-to-cblinfun-code*[*code*]:
— Code equation for translating a vector into an operation (single-column matrix)
 $mat\text{-of-cblinfun } (vector\text{-to-cblinfun-code } \psi) = mat\text{-of-cols } (CARD('a)) \ [vec\text{-of-ell2 } \psi]$
for $\psi :: 'a :: enum \ ell2$
 $\langle proof \rangle$

definition *butterfly-code* :: '<*a ell2* $\Rightarrow$ '*b ell2* $\Rightarrow$ '*b ell2* $\Rightarrow_{CL}$ '*a ell2*>

where [*code del*, *code-abbrev*]: $\langle butterfly\text{-code} = butterfly \rangle$

lemma *butterfly-code*[*code*]: $\langle mat\text{-of-cblinfun } (butterfly\text{-code } s \ t)$
 $= mat\text{-of-cols } (CARD('a)) \ [vec\text{-of-ell2 } s] * mat\text{-of-rows } (CARD('b)) \ [map\text{-vec } cnj \ (vec\text{-of-ell2 } t)] \rangle$
for $s :: 'a :: enum \ ell2$ **and** $t :: 'b :: enum \ ell2$
 $\langle proof \rangle$

17.4 Subspaces

In this section, we define code equations for handling subspaces, i.e., values of type '*a ccspace*. We choose to computationally represent a subspace by a list of vectors that span the subspace. That is, if *vecs* are vectors

(type *complex vec*), *SPAN vecs* is defined to be their span. Then the code generation can simply represent all subspaces in this form, and we need to define the operations on subspaces in terms of list of vectors (e.g., the closed union of two subspaces would be computed as the concatenation of the two lists, to give one of the simplest examples).

To support this, *SPAN* is declared as a "code-datatype". (Not as an abstract datatype like *cblinfun-of-mat/mat-of-cblinfun* because that would require *SPAN* to be injective.)

Then all code equations for different operations need to be formulated as functions of values of the form *SPAN x*. (E.g., *SPAN x + SPAN y = SPAN (...)*.)

definition [code del]: *SPAN x = (let n = length (canonical-basis :: 'a::onb-enum list) in*

ccspan (basis-enum-of-vec 'Set.filter (λv. dim-vec v = n) (set x)) :: 'a ccspace)

— The *SPAN* of vectors *x*, as a *ccspace*. We filter out vectors of the wrong dimension because *SPAN* needs to have well-defined behavior even in cases that would not actually occur in an execution.

code-datatype *SPAN*

We first declare code equations for *Proj*, i.e., for turning a subspace into a projector. This means, we would need a code equation of the form *mat-of-cblinfun (Proj (SPAN S)) = ...*. However, this equation is not accepted by the code generation for reasons we do not understand. But if we define an auxiliary constant *mat-of-cblinfun-Proj-code* that stands for *mat-of-cblinfun (Proj -)*, define a code equation for *mat-of-cblinfun-Proj-code*, and then define a code equation for *mat-of-cblinfun (Proj S)* in terms of *mat-of-cblinfun-Proj-code*, Isabelle accepts the code equations.

definition *mat-of-cblinfun-Proj-code S = mat-of-cblinfun (Proj S)*

declare *mat-of-cblinfun-Proj-code-def[symmetric, code]*

lemma *mat-of-cblinfun-Proj-code-code[code]*:

— Code equation for computing a projector onto a set *S* of vectors. We first make the vectors *S* into an orthonormal basis using the Gram-Schmidt procedure and then compute the projector as the sum of the "butterflies" *x * x** of the vectors *x ∈ S* (done by *mk-projector*).

mat-of-cblinfun-Proj-code (SPAN S :: 'a::onb-enum ccspace) =
(let d = length (canonical-basis :: 'a list) in mk-projector d (filter (λv. dim-vec v = d) S))
<proof>

lemma *top-ccspace-code[code]*:

— Code equation for $\top$, the subspace containing everything. Top is represented as the span of the standard basis vectors.

(top::'a ccspace) =
(let n = length (canonical-basis :: 'a::onb-enum list) in SPAN (unit-vecs n))
<proof>

lemma *bot-as-span*[code]:

— Code equation for $\perp$, the subspace containing everything. Top is represented as the span of the standard basis vectors.

(*bot*::'a::onb-enum *ccsubspace*) = *SPAN* []

⟨*proof*⟩

lemma *sup-spans*[code]:

— Code equation for the join (lub) of two subspaces (union of the generating lists)

SPAN A $\sqcup$ *SPAN* B = *SPAN* (A @ B)

⟨*proof*⟩

We do not need an equation for (+) because (+) is defined in terms of ($\sqcup$) (for *ccsubspace*), thus the code generation automatically computes (+) in terms of the code for ($\sqcup$)

definition [code del,code-abbrev]: *Span-code* (*S*::'a::enum *ell2* set) = (*ccspan* *S*)

— A copy of *ccspan* with restricted type. For analogous reasons as *cblin-fun-apply-ell2*, see there for explanations

lemma *span-Set-Monad*[code]: *Span-code* (*Set-Monad* l) = (*SPAN* (map *vec-of-ell2* l))

— Code equation for the span of a finite set. (*Set-Monad* is a datatype constructor that represents sets as lists in the computation.)

⟨*proof*⟩

This instantiation defines a code equation for equality tests for *ccsubspace*. The actual code for equality tests is given below (lemma *equal-ccsubspace-code*).

instantiation *ccsubspace* :: (onb-enum) equal **begin**

definition [code del]: *equal-ccsubspace* (*A*::'a *ccsubspace*) *B* = (*A=B*)

instance ⟨*proof*⟩

end

lemma *leq-ccsubspace-code*[code]:

— Code equation for deciding inclusion of one space in another. Uses the constant *is-subspace-of-vec-list* which implements the actual computation by checking for each generator of A whether it is in the span of B (by orthogonal projection onto an orthonormal basis of B which is computed using Gram-Schmidt).

SPAN A $\leq$ (*SPAN* B :: 'a::onb-enum *ccsubspace*)

$\longleftrightarrow$ (let *d* = length (canonical-basis :: 'a list) in

is-subspace-of-vec-list *d*

(filter ($\lambda v. \text{dim-vec } v = d$) *A*)

(filter ($\lambda v. \text{dim-vec } v = d$) *B*))

⟨*proof*⟩

lemma *equal-ccsubspace-code*[code]:

— Code equation for equality test. By checking mutual inclusion (for which we have code by the preceding code equation).

HOL.equal (*A* ::- *ccsubspace*) *B* = (*A* ≤ *B* ∧ *B* ≤ *A*)
 ⟨*proof*⟩

lemma *cblinfun-image-code*[*code*]:

— Code equation for applying an operator *A* to a subspace. Simply by multiplying each generator with *A*

A *_{*S*} *SPAN S* = (let *d* = length (canonical-basis :: 'a list) in
 SPAN (map (mult-mat-vec (mat-of-cblinfun *A*))
 (filter (λ*v*. dim-vec *v* = *d*) *S*)))

for *A* :: 'a :: onb-enum ⇒_{*CL*} 'b :: onb-enum

⟨*proof*⟩

definition [*code del*, *code-abbrev*]: *range-cblinfun-code* *A* = *A* *_{*S*} *top*

— A new constant for the special case of applying an operator to the subspace *T* (i.e., for computing the range of the operator). We do this to be able to give more specialized code for this specific situation. (The generic code for (*_{*S*}) would work but is less efficient because it involves repeated matrix multiplications. *code-abbrev* makes sure occurrences of *A* *_{*S*} *T* are replaced before starting the actual code generation.

lemma *range-cblinfun-code*[*code*]:

— Code equation for computing the range of an operator *A*. Returns the columns of the matrix representation of *A*.

fixes *A* :: 'a :: onb-enum ⇒_{*CL*} 'b :: onb-enum

shows *range-cblinfun-code* *A* = *SPAN* (*cols* (mat-of-cblinfun *A*))

⟨*proof*⟩

lemma *uminus-Span-code*[*code*]: — *X* = *range-cblinfun-code* (*id-cblinfun* — *Proj X*)

— Code equation for the orthogonal complement of a subspace *X*. Computed as the range of one minus the projector on *X*

⟨*proof*⟩

lemma *kernel-code*[*code*]:

— Computes the kernel of an operator *A*. This is implemented using the existing functions for transforming a matrix into row echelon form (*gauss-jordan-single*) and for computing a basis of the kernel of such a matrix (*find-base-vectors*)

kernel *A* = *SPAN* (*find-base-vectors* (*gauss-jordan-single* (mat-of-cblinfun *A*)))

for *A* :: ('a :: onb-enum, 'b :: onb-enum) cblinfun

⟨*proof*⟩

lemma *inf-ccsubspace-code*[*code*]:

— Code equation for intersection of subspaces. Reduced to orthogonal complement and sum of subspaces for which we already have code equations.

(*A* :: 'a :: onb-enum ccsubspace) □ *B* = — (— *A* □ — *B*)

⟨*proof*⟩

lemma *Sup-ccsubspace-code*[*code*]:

— Supremum (sum) of a set of subspaces. Implemented by repeated pairwise sum.

```
Sup (Set-Monad l :: 'a::onb-enum ccspace set) = fold sup l bot
⟨proof⟩
```

lemma *Inf-ccspace-code*[code]:

— Infimum (intersection) of a set of subspaces. Implemented by the orthogonal complement of the supremum.

```
Inf (Set-Monad l :: 'a::onb-enum ccspace set)
= - Sup (Set-Monad (map uminus l))
⟨proof⟩
```

17.5 Miscellanea

This is a hack to circumvent a bug in the code generation. The automatically generated code for the class *uniformity* has a type that is different from what the generated code later assumes, leading to compilation errors (in ML at least) in any expression involving *- ell2* (even if the constant *uniformity* is not actually used).

The fragment below circumvents this by forcing Isabelle to use the right type. (The logically useless fragment "*let x = ((=)::'a⇒⇒-)*" achieves this.)

```
lemma uniformity-ell2-code[code]: (uniformity :: ('a ell2 * -) filter) = Filter.abstract-filter
(%-.
  Code.abort STR "no uniformity" (%-.
    let x = ((=)::'a⇒⇒-) in uniformity))
⟨proof⟩
```

Code equation for *UNIV*. It is now implemented via type class *enum* (which provides a list of all values).

```
declare enum-class.UNIV-enum[code]
```

Setup for code generation involving sets of *ell2/ccspace*. This configures to use lists for representing sets in code.

```
derive (eq) ceq ccspace
derive (no) ccompare ccspace
derive (monad) set-impl ccspace
derive (eq) ceq ell2
derive (no) ccompare ell2
derive (monad) set-impl ell2
```

```
unbundle no lattice-syntax and no jnf-syntax and no cblinfun-syntax
```

```
end
```

18 *Cblinfun-Code-Examples* – Examples and test cases for code generation

```
theory Cblinfun-Code-Examples
imports
  Complex-Bounded-Operators.Extra-Pretty-Code-Examples
  Jordan-Normal-Form.Matrix-Impl
  HOL-Library.Code-Target-Numeral
  Cblinfun-Code
begin

hide-const (open) Order.bottom Order.top
no-notation Lattice.join (infixl  $\langle \sqcup_1 \rangle$  65)
no-notation Lattice.meet (infixl  $\langle \sqcap_1 \rangle$  70)

unbundle lattice-syntax
unbundle cblinfun-syntax
```

19 Examples

19.1 Operators

```
value id-cblinfun :: bool ell2  $\Rightarrow_{CL}$  bool ell2

value 1 :: unit ell2  $\Rightarrow_{CL}$  unit ell2

value id-cblinfun + id-cblinfun :: bool ell2  $\Rightarrow_{CL}$  bool ell2

value 0 :: (bool ell2  $\Rightarrow_{CL}$  Enum.finite-3 ell2)

value  $-$  id-cblinfun :: bool ell2  $\Rightarrow_{CL}$  bool ell2

value id-cblinfun  $-$  id-cblinfun :: bool ell2  $\Rightarrow_{CL}$  bool ell2

value classical-operator ( $\lambda b.$  Some ( $\neg$  b))

value  $\langle$  explicit-cblinfun ( $\lambda x y ::$  bool. of-bool ( $x \wedge y$ ))  $\rangle$ 

value id-cblinfun = (0 :: bool ell2  $\Rightarrow_{CL}$  bool ell2)

value 2 *R id-cblinfun :: bool ell2  $\Rightarrow_{CL}$  bool ell2

value imaginary-unit *C id-cblinfun :: bool ell2  $\Rightarrow_{CL}$  bool ell2

value id-cblinfun oCL 0 :: bool ell2  $\Rightarrow_{CL}$  bool ell2

value id-cblinfun* :: bool ell2  $\Rightarrow_{CL}$  bool ell2
```

19.2 Vectors

value $0 :: \text{bool ell2}$

value $1 :: \text{unit ell2}$

value ket False

value $2 *_C \text{ket False}$

value $2 *_R \text{ket False}$

value $\text{ket True} + \text{ket False}$

value $\text{ket True} - \text{ket True}$

value $\text{ket True} = \text{ket True}$

value $- \text{ket True}$

value $\text{cinner} (\text{ket True}) (\text{ket True})$

value $\text{norm} (\text{ket True})$

value $\text{ket} () * \text{ket} ()$

value $1 :: \text{unit ell2}$

value $(1::\text{unit ell2}) * (1::\text{unit ell2})$

19.3 Vector/Matrix

value $\text{id-cblinfun} *_V \text{ket True}$

value $\langle \text{vector-to-cblinfun} (\text{ket True}) :: \text{unit ell2} \Rightarrow_{CL} \rightarrow$

19.4 Subspaces

value $\text{ccspan} \{\text{ket False}\}$

value $\text{Proj} (\text{ccspan} \{\text{ket False}\})$

value $\text{top} :: \text{bool ell2 ccsubspace}$

value $\text{bot} :: \text{bool ell2 ccsubspace}$

value $0 :: \text{bool ell2 ccsubspace}$

value $\text{ccspan} \{\text{ket False}\} \sqcup \text{ccspan} \{\text{ket True}\}$


```

value ccspan {ket False} + ccspan {ket True}

value ccspan {ket False}  $\sqcap$  ccspan {ket True}

value id-cblinfun *S ccspan {ket False}

value id-cblinfun *S (top :: bool ell2 ccspace)

value − ccspan {ket False}

value ccspan {ket False, ket True} = top

value ccspan {ket False} ≤ ccspan {ket True}

value cblinfun-image id-cblinfun (ccspan {ket True})

value kernel id-cblinfun :: bool ell2 ccspace

value eigenspace 1 id-cblinfun :: bool ell2 ccspace

value Inf {ccspan {ket False}, top}

value Sup {ccspan {ket False}, top}

end

```

References

- [1] J. B. Conway. *A course in functional analysis*, volume 96. Springer Science & Business Media, 2013.
- [2] F. Haftmann. Code generation from Isabelle/HOL theories. <https://isabelle.in.tum.de/website-Isabelle2019/dist/Isabelle2019/doc/codegen.pdf>, 2019.