

Operations on Bounded Natural Functors

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Abstract

This entry formalizes the closure property of bounded natural functors (BNFs) under seven operations. These operations and the corresponding proofs constitute the core of Isabelle’s (co)datatype package. To be close to the implemented tactics, the proofs are deliberately formulated as detailed apply scripts. The (co)datatypes together with (co)induction principles and (co)recursors are byproducts of the fixpoint operations LFP and GFP. Composition of BNFs is subdivided into four simpler operations: Compose, Kill, Lift, and Permute. The N2M operation provides mutual (co)induction principles and (co)recursors for nested (co)datatypes.

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1 Least Fixpoint (a.k.a. Datatype)

$\langle ML \rangle$

notation *BNF_Def.convol* ($\langle _ , _ \rangle$)

'b1 = ('a, 'b1, 'b2) F1

'b2 = ('a, 'b1, 'b2) F2

To build a witness scenario, let us assume

('a, 'b1, 'b2) F1 = 'a * 'b1 + 'a * 'b2

('a, 'b1, 'b2) F2 = unit + 'b1 * 'b2

declare *[[bnf_internals]]*

bnf-axiomatization (*F1set1: 'a, F1set2: 'b1, F1set3: 'b2*) *F1*

[wits: 'a \Rightarrow 'b1 \Rightarrow ('a, 'b1, 'b2) F1 'a \Rightarrow 'b2 \Rightarrow ('a, 'b1, 'b2) F1]

for *map: F1map rel: F1rel*

bnf-axiomatization (*F2set1: 'a, F2set2: 'b1, F2set3: 'b2*) *F2*

[wits: ('a, 'b1, 'b2) F2]

for *map: F2map rel: F2rel*

abbreviation *F1in :: 'a1 set \Rightarrow 'a2 set \Rightarrow 'a3 set \Rightarrow (('a1, 'a2, 'a3) F1) set* **where**

F1in A1 A2 A3 \equiv {x. F1set1 x \subseteq A1 \wedge F1set2 x \subseteq A2 \wedge F1set3 x \subseteq A3}

abbreviation *F2in :: 'a1 set \Rightarrow 'a2 set \Rightarrow 'a3 set \Rightarrow (('a1, 'a2, 'a3) F2) set* **where**

F2in A1 A2 A3 \equiv {x. F2set1 x \subseteq A1 \wedge F2set2 x \subseteq A2 \wedge F2set3 x \subseteq A3}

lemma *F1map_comp_id: F1map g1 g2 g3 (F1map id f2 f3 x) = F1map g1 (g2 o f2) (g3 o f3) x*

\langle proof \rangle

lemmas *F1in_mono23 = F1.in_mono[OF subset_refl]*

lemma *F1map_congL: $\llbracket \forall a \in F1set2 x. f a = a; \forall a \in F1set3 x. g a = a \rrbracket \Longrightarrow$*

F1map id f g x = x

\langle proof \rangle

lemma *F2map_comp_id: F2map g1 g2 g3 (F2map id f2 f3 x) = F2map g1 (g2 o f2) (g3 o f3) x*

\langle proof \rangle

lemmas *F2in_mono23 = F2.in_mono[OF subset_refl]*

lemma *F2map_congL: $\llbracket \forall a \in F2set2 x. f a = a; \forall a \in F2set3 x. g a = a \rrbracket \Longrightarrow$*

F2map id f g x = x

\langle proof \rangle

1.1 Algebra

definition *alg* **where**

alg B1 B2 s1 s2 =

(($\forall x \in F1in$ (UNIV :: 'a set) B1 B2. s1 x \in B1) \wedge ($\forall y \in F2in$ (UNIV :: 'a set) B1 B2. s2 y \in B2))

lemma *alg_F1set: $\llbracket alg B1 B2 s1 s2; F1set2 x \subseteq B1; F1set3 x \subseteq B2 \rrbracket \Longrightarrow s1 x \in B1$*

\langle proof \rangle

lemma *alg_F2set: $\llbracket alg B1 B2 s1 s2; F2set2 x \subseteq B1; F2set3 x \subseteq B2 \rrbracket \Longrightarrow s2 x \in B2$*

<proof>

lemma *alg_not_empty*:

$alg\ B1\ B2\ s1\ s2 \implies B1 \neq \{\} \wedge B2 \neq \{\}$

<proof>

1.2 Morphism

definition *mor* **where**

$mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f\ g =$
 $((\forall a \in B1. f\ a \in B1') \wedge (\forall a \in B2. g\ a \in B2')) \wedge$
 $((\forall z \in F1in\ (UNIV :: 'a\ set)\ B1\ B2. f\ (s1\ z) = s1'\ (F1map\ id\ f\ g\ z)) \wedge$
 $(\forall z \in F2in\ (UNIV :: 'a\ set)\ B1\ B2. g\ (s2\ z) = s2'\ (F2map\ id\ f\ g\ z)))$

lemma *morE1*: $\llbracket mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f\ g; z \in F1in\ UNIV\ B1\ B2 \rrbracket$

$\implies f\ (s1\ z) = s1'\ (F1map\ id\ f\ g\ z)$

<proof>

lemma *morE2*: $\llbracket mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f\ g; z \in F2in\ UNIV\ B1\ B2 \rrbracket$

$\implies g\ (s2\ z) = s2'\ (F2map\ id\ f\ g\ z)$

<proof>

lemma *mor_incl*: $\llbracket B1 \subseteq B1'; B2 \subseteq B2' \rrbracket \implies mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1\ s2\ id\ id$

<proof>

lemma *mor_comp*:

$\llbracket mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f\ g;$
 $mor\ B1'\ B2'\ s1'\ s2'\ B1''\ B2''\ s1''\ s2''\ f'\ g' \rrbracket \implies$
 $mor\ B1\ B2\ s1\ s2\ B1''\ B2''\ s1''\ s2''\ (f' \circ f)\ (g' \circ g)$

<proof>

lemma *mor_cong*: $\llbracket f' = f; g' = g; mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f\ g \rrbracket \implies$

$mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f'\ g'$

<proof>

lemma *mor_str*:

$mor\ UNIV\ UNIV\ (F1map\ id\ s1\ s2)\ (F2map\ id\ s1\ s2)\ UNIV\ UNIV\ s1\ s2\ s1\ s2$

<proof>

1.3 Bounds

type-synonym $bd_type_F1' = bd_type_F1 + (bd_type_F1, bd_type_F1, bd_type_F1)\ F1$

type-synonym $bd_type_F2' = bd_type_F2 + (bd_type_F2, bd_type_F2, bd_type_F2)\ F2$

type-synonym $SucFbd_type = ((bd_type_F1' + bd_type_F2')\ set)$

type-synonym $'a1\ ASucFbd_type = (SucFbd_type \Rightarrow ('a1 + bool))$

abbreviation $F1bd' \equiv bd_F1 + c \mid UNIV :: (bd_type_F1, bd_type_F1, bd_type_F1)\ F1\ set$

lemma *F1set1_bd_incr*: $\bigwedge x. |F1set1\ x| \leq_o\ F1bd'$

<proof>

lemma *F1set2_bd_incr*: $\bigwedge x. |F1set2\ x| \leq_o\ F1bd'$

<proof>

lemma *F1set3_bd_incr*: $\bigwedge x. |F1set3\ x| \leq_o\ F1bd'$

<proof>

lemmas $F1bd'_Card_order = Card_order_csum$

lemmas $F1bd'_Cinfinite = Cinfinite_csum1[OF\ F1.bd_Cinfinite]$

lemmas $F1bd'_Cnotzero = Cinfinite_Cnotzero[OF\ F1bd'_Cinfinite]$

lemmas $F1bd'_card_order = card_order_csum[OF\ F1.bd_card_order_card_of_card_order_on]$

abbreviation $F2bd' \equiv bd_F2 + c \mid UNIV :: (bd_type_F2, bd_type_F2, bd_type_F2)\ F2\ set$

lemma *F2set1_bd_incr*: $\bigwedge x. |F2set1\ x| \leq_o\ F2bd'$

<proof>

lemma *F2set2_bd_incr*: $\bigwedge x. |F2set2\ x| \leq_o\ F2bd'$

<proof>

lemma $F2set3_bd_incr$: $\bigwedge x. |F2set3\ x| \leq_o F2bd'$
 ⟨proof⟩

lemmas $F2bd'_Card_order = Card_order_csum$

lemmas $F2bd'_Cinfinite = Cinfinite_csum1[OF\ F2.bd_Cinfinite]$

lemmas $F2bd'_Cnotzero = Cinfinite_Cnotzero[OF\ F2bd'_Cinfinite]$

lemmas $F2bd'_card_order = card_order_csum[OF\ F2.bd_card_order\ card_of_card_order_on]$

abbreviation $SucFbd$ **where** $SucFbd \equiv cardSuc\ (F1bd' +_c\ F2bd')$

abbreviation $ASucFbd$ **where** $ASucFbd \equiv (|UNIV| +_c\ ctwo) \wedge_c\ SucFbd$

lemma $F1set1_bd'$: $|F1set1\ x| \leq_o F1bd' +_c\ F2bd'$
 ⟨proof⟩

lemma $F1set2_bd'$: $|F1set2\ x| \leq_o F1bd' +_c\ F2bd'$
 ⟨proof⟩

lemma $F1set3_bd'$: $|F1set3\ x| \leq_o F1bd' +_c\ F2bd'$
 ⟨proof⟩

lemma $F2set1_bd'$: $|F2set1\ x| \leq_o F1bd' +_c\ F2bd'$
 ⟨proof⟩

lemma $F2set2_bd'$: $|F2set2\ x| \leq_o F1bd' +_c\ F2bd'$
 ⟨proof⟩

lemma $F2set3_bd'$: $|F2set3\ x| \leq_o F1bd' +_c\ F2bd'$
 ⟨proof⟩

lemmas $SucFbd_Card_order = cardSuc_Card_order[OF\ Card_order_csum]$

lemmas $SucFbd_Cinfinite = Cinfinite_cardSuc[OF\ Cinfinite_csum1[OF\ F1bd'_Cinfinite]]$

lemmas $SucFbd_Cnotzero = Cinfinite_Cnotzero[OF\ SucFbd_Cinfinite]$

lemmas $worel_SucFbd = Card_order_wo_rel[OF\ SucFbd_Card_order]$

lemmas $ASucFbd_Cinfinite = Cinfinite_cexp[OF\ ordLeq_csum2[OF\ Card_order_ctwo]\ SucFbd_Cinfinite]$

1.4 Minimal Algebras

abbreviation min_G1 **where**

$min_G1\ As1_As2\ i \equiv (\bigcup j \in underS\ SucFbd\ i.\ fst\ (As1_As2\ j))$

abbreviation min_G2 **where**

$min_G2\ As1_As2\ i \equiv (\bigcup j \in underS\ SucFbd\ i.\ snd\ (As1_As2\ j))$

abbreviation min_H **where**

$min_H\ s1\ s2\ As1_As2\ i \equiv$

$(min_G1\ As1_As2\ i \cup s1 \text{ ' } (F1in\ (UNIV :: 'a\ set)\ (min_G1\ As1_As2\ i)\ (min_G2\ As1_As2\ i))),$
 $min_G2\ As1_As2\ i \cup s2 \text{ ' } (F2in\ (UNIV :: 'a\ set)\ (min_G1\ As1_As2\ i)\ (min_G2\ As1_As2\ i)))$

abbreviation min_algs **where**

$min_algs\ s1\ s2 \equiv wo_rel.worec\ SucFbd\ (min_H\ s1\ s2)$

definition min_alg1 **where**

$min_alg1\ s1\ s2 = (\bigcup i \in Field\ SucFbd.\ fst\ (min_algs\ s1\ s2\ i))$

definition min_alg2 **where**

$min_alg2\ s1\ s2 = (\bigcup i \in Field\ SucFbd.\ snd\ (min_algs\ s1\ s2\ i))$

lemma min_algs :

$i \in Field\ SucFbd \implies min_algs\ s1\ s2\ i = min_H\ s1\ s2\ (min_algs\ s1\ s2)\ i$

⟨proof⟩

corollary min_algs1 : $i \in Field\ SucFbd \implies fst\ (min_algs\ s1\ s2\ i) =$

$min_G1\ (min_algs\ s1\ s2)\ i \cup$

$s1 \text{ ' } (F1in\ UNIV\ (min_G1\ (min_algs\ s1\ s2)\ i)\ (min_G2\ (min_algs\ s1\ s2)\ i))$

<proof>

corollary $\text{min_algs2}: i \in \text{Field SucFbd} \implies \text{snd} (\text{min_algs } s1 \ s2 \ i) =$
 $\text{min_G2} (\text{min_algs } s1 \ s2) \ i \cup$
 $s2 \text{' } (\text{F2in UNIV} (\text{min_G1} (\text{min_algs } s1 \ s2) \ i) (\text{min_G2} (\text{min_algs } s1 \ s2) \ i))$
<proof>

lemma $\text{min_algs_mono1}: \text{relChain SucFbd } (\%i. \text{fst} (\text{min_algs } s1 \ s2 \ i))$
<proof>

lemma $\text{min_algs_mono2}: \text{relChain SucFbd } (\%i. \text{snd} (\text{min_algs } s1 \ s2 \ i))$
<proof>

lemma $\text{SucFbd_limit}: \llbracket x1 \in \text{Field SucFbd} \ \& \ x2 \in \text{Field SucFbd} \rrbracket$
 $\implies \exists y \in \text{Field SucFbd}. (x1 \neq y \wedge (x1, y) \in \text{SucFbd}) \wedge (x2 \neq y \wedge (x2, y) \in \text{SucFbd})$
<proof>

lemma $\text{alg_min_alg}: \text{alg} (\text{min_alg1 } s1 \ s2) (\text{min_alg2 } s1 \ s2) \ s1 \ s2$
<proof>

lemmas $\text{SucFbd_ASucFbd} = \text{ordLess_ordLeq_trans}[OF$
 ordLess_ctwo_cexp
 $\text{cexp_mono1}[OF \text{ordLeq_csum2}[OF \text{Card_order_ctwo}]],$
 $OF \text{SucFbd_Card_order SucFbd_Card_order}]$

lemma card_of_min_algs :
fixes $s1 :: ('a, 'b, 'c) F1 \Rightarrow 'b$ **and** $s2 :: ('a, 'b, 'c) F2 \Rightarrow 'c$
shows $i \in \text{Field SucFbd} \longrightarrow$
 $(|\text{fst} (\text{min_algs } s1 \ s2 \ i)| \leq_o (\text{ASucFbd} :: 'a \ \text{ASucFbd_type rel}) \wedge |\text{snd} (\text{min_algs } s1 \ s2 \ i)| \leq_o (\text{ASucFbd} :: 'a \ \text{ASucFbd_type rel}))$
<proof>

lemma card_of_min_alg1 :
fixes $s1 :: ('a, 'b, 'c) F1 \Rightarrow 'b$ **and** $s2 :: ('a, 'b, 'c) F2 \Rightarrow 'c$
shows $|\text{min_alg1 } s1 \ s2| \leq_o (\text{ASucFbd} :: 'a \ \text{ASucFbd_type rel})$
<proof>

lemma card_of_min_alg2 :
fixes $s1 :: ('a, 'b, 'c) F1 \Rightarrow 'b$ **and** $s2 :: ('a, 'b, 'c) F2 \Rightarrow 'c$
shows $|\text{min_alg2 } s1 \ s2| \leq_o (\text{ASucFbd} :: 'a \ \text{ASucFbd_type rel})$
<proof>

lemma $\text{least_min_algs}: \text{alg } B1 \ B2 \ s1 \ s2 \implies$
 $i \in \text{Field SucFbd} \longrightarrow$
 $\text{fst} (\text{min_algs } s1 \ s2 \ i) \subseteq B1 \wedge \text{snd} (\text{min_algs } s1 \ s2 \ i) \subseteq B2$
<proof>

lemma $\text{least_min_alg1}: \text{alg } B1 \ B2 \ s1 \ s2 \implies \text{min_alg1 } s1 \ s2 \subseteq B1$
<proof>

lemma $\text{least_min_alg2}: \text{alg } B1 \ B2 \ s1 \ s2 \implies \text{min_alg2 } s1 \ s2 \subseteq B2$
<proof>

lemma mor_incl_min_alg :
 $\text{alg } B1 \ B2 \ s1 \ s2 \implies$
 $\text{mor} (\text{min_alg1 } s1 \ s2) (\text{min_alg2 } s1 \ s2) \ s1 \ s2 \ B1 \ B2 \ s1 \ s2 \ \text{id} \ \text{id}$
<proof>

1.5 Initiality

The following “happens” to be the type (for our particular construction) of the initial algebra carrier:

type-synonym $'a1 \ F1\text{init_type} = ('a1, 'a1 \ \text{ASucFbd_type}, 'a1 \ \text{ASucFbd_type}) \ F1 \Rightarrow 'a1 \ \text{ASucFbd_type}$
type-synonym $'a1 \ F2\text{init_type} = ('a1, 'a1 \ \text{ASucFbd_type}, 'a1 \ \text{ASucFbd_type}) \ F2 \Rightarrow 'a1 \ \text{ASucFbd_type}$

typedef 'a1 IIT =
 UNIV ::
 (('a1 ASucFbd_type set × 'a1 ASucFbd_type set) × ('a1 F1init_type × 'a1 F2init_type)) set
 ⟨proof⟩

1.6 Initial Algebras

abbreviation II :: 'a1 IIT set **where**

II ≡ {Abs_IIT ((B1, B2), (s1, s2)) | B1 B2 s1 s2. alg B1 B2 s1 s2}

definition str_init1 **where**

str_init1 (dummy :: 'a1)
 (y :: ('a1, 'a1 IIT ⇒ 'a1 ASucFbd_type, 'a1 IIT ⇒ 'a1 ASucFbd_type) F1)
 (i :: 'a1 IIT) =
 fst (snd (Rep_IIT i))
 (F1map id (λf :: 'a1 IIT ⇒ 'a1 ASucFbd_type. f i) (λf. f i) y)

definition str_init2 **where**

str_init2 (dummy :: 'a1) y (i :: 'a1 IIT) =
 snd (snd (Rep_IIT i)) (F2map id (λf. f i) (λf. f i) y)

abbreviation car_init1 **where**

car_init1 dummy ≡ min_alg1 (str_init1 dummy) (str_init2 dummy)

abbreviation car_init2 **where**

car_init2 dummy ≡ min_alg2 (str_init1 dummy) (str_init2 dummy)

lemma alg_select:

∀ i ∈ II. alg (fst (fst (Rep_IIT i))) (snd (fst (Rep_IIT i)))
 (fst (snd (Rep_IIT i))) (snd (snd (Rep_IIT i)))
 ⟨proof⟩

lemma mor_select:

[[i ∈ II;
 mor (fst (fst (Rep_IIT i))) (snd (fst (Rep_IIT i)))
 (fst (snd (Rep_IIT i))) (snd (snd (Rep_IIT i))) UNIV UNIV s1' s2' f g]] ⇒
 mor (car_init1 dummy) (car_init2 dummy) (str_init1 dummy) (str_init2 dummy) UNIV UNIV s1' s2' (f ∘ (λh.
 h i)) (g ∘ (λh. h i))
 ⟨proof⟩

lemma init_unique_mor:

[[a1 ∈ car_init1 dummy; a2 ∈ car_init2 dummy;
 mor (car_init1 dummy) (car_init2 dummy) (str_init1 dummy) (str_init2 dummy) B1 B2 s1 s2 f1 f2;
 mor (car_init1 dummy) (car_init2 dummy) (str_init1 dummy) (str_init2 dummy) B1 B2 s1 s2 g1 g2]] ⇒
 f1 a1 = g1 a1 ∧ f2 a2 = g2 a2
 ⟨proof⟩

abbreviation closed **where**

closed dummy phi1 phi2 ≡ ((∀ x ∈ F1in UNIV (car_init1 dummy) (car_init2 dummy).
 (∀ z ∈ F1set2 x. phi1 z) ∧ (∀ z ∈ F1set3 x. phi2 z) → phi1 (str_init1 dummy x)) ∧
 (∀ x ∈ F2in UNIV (car_init1 dummy) (car_init2 dummy).
 (∀ z ∈ F2set2 x. phi1 z) ∧ (∀ z ∈ F2set3 x. phi2 z) → phi2 (str_init2 dummy x)))

lemma init_induct: closed dummy phi1 phi2 ⇒

(∀ x ∈ car_init1 dummy. phi1 x) ∧ (∀ x ∈ car_init2 dummy. phi2 x)
 ⟨proof⟩

1.7 The datatype

typedef (overloaded) 'a1 IF1 = car_init1 (undefined :: 'a1)
 ⟨proof⟩

typedef (overloaded) 'a1 IF2 = car_init2 (undefined :: 'a1)
 ⟨proof⟩

definition ctor1 **where** ctor1 = Abs_IF1 o str_init1 undefined o F1map id Rep_IF1 Rep_IF2

definition ctor2 **where** ctor2 = Abs_IF2 o str_init2 undefined o F2map id Rep_IF1 Rep_IF2

lemma *mor_Rep_IF*:

mor (*UNIV* :: 'a *IF1* set) (*UNIV* :: 'a *IF2* set) *ctor1* *ctor2*
(*car_init1* undefined) (*car_init2* undefined) (*str_init1* undefined) (*str_init2* undefined) *Rep_IF1* *Rep_IF2*
<proof>

lemma *mor_Abs_IF*:

mor (*car_init1* undefined) (*car_init2* undefined)
(*str_init1* undefined) (*str_init2* undefined) *UNIV* *UNIV* *ctor1* *ctor2* *Abs_IF1* *Abs_IF2*
<proof>

lemma *copy*:

$\llbracket \text{alg } B1 \ B2 \ s1 \ s2; \text{bij_betw } f \ B1' \ B1; \text{bij_betw } g \ B2' \ B2 \rrbracket \implies$
 $\exists f' \ g'. \text{alg } B1' \ B2' \ f' \ g' \wedge \text{mor } B1' \ B2' \ f' \ g' \ B1 \ B2 \ s1 \ s2 \ f \ g$
<proof>

lemma *init_ex_mor*:

$\exists f \ g. \text{mor } UNIV \ UNIV \ \text{ctor1} \ \text{ctor2} \ UNIV \ UNIV \ s1 \ s2 \ f \ g$
<proof>

Iteration

abbreviation *fold* **where**

fold *s1* *s2* \equiv (*SOME* *f*. *mor* *UNIV* *UNIV* *ctor1* *ctor2* *UNIV* *UNIV* *s1* *s2* (*fst* *f*) (*snd* *f*))

definition *fold1* **where** *fold1* *s1* *s2* = *fst* (*fold* *s1* *s2*)

definition *fold2* **where** *fold2* *s1* *s2* = *snd* (*fold* *s1* *s2*)

lemma *mor_fold*:

mor *UNIV* *UNIV* *ctor1* *ctor2* *UNIV* *UNIV* *s1* *s2* (*fold1* *s1* *s2*) (*fold2* *s1* *s2*)
<proof>

<ML>

theorem *fold1*:

(*fold1* *s1* *s2*) (*ctor1* *x*) = *s1* (*F1map* *id* (*fold1* *s1* *s2*) (*fold2* *s1* *s2*) *x*)
<proof>

theorem *fold2*:

(*fold2* *s1* *s2*) (*ctor2* *x*) = *s2* (*F2map* *id* (*fold1* *s1* *s2*) (*fold2* *s1* *s2*) *x*)
<proof>

lemma *mor_UNIV*: *mor* *UNIV* *UNIV* *s1* *s2* *UNIV* *UNIV* *s1'* *s2'* *f* *g* \longleftrightarrow

f *o* *s1* = *s1'* *o* *F1map* *id* *f* *g* \wedge *g* *o* *s2* = *s2'* *o* *F2map* *id* *f* *g*
<proof>

lemma *fold_unique_mor*: *mor* *UNIV* *UNIV* *ctor1* *ctor2* *UNIV* *UNIV* *s1* *s2* *f* *g* \implies

f = *fold1* *s1* *s2* \wedge *g* = *fold2* *s1* *s2*
<proof>

lemmas *fold_unique* = *fold_unique_mor*[*OF* *iffD2*[*OF* *mor_UNIV*], *OF* *conjI*]

lemmas *fold1_ctor* = *sym*[*OF* *conjunct1*[*OF* *fold_unique_mor*[*OF* *mor_incl*[*OF* *subset_UNIV* *subset_UNIV*]]]]

lemmas *fold2_ctor* = *sym*[*OF* *conjunct2*[*OF* *fold_unique_mor*[*OF* *mor_incl*[*OF* *subset_UNIV* *subset_UNIV*]]]]

Case distinction

lemmas *ctor1_o_fold1* =

trans[*OF* *conjunct1*[*OF* *fold_unique_mor*[*OF* *mor_comp*[*OF* *mor_fold* *mor_str*]]]] *fold1_ctor*

lemmas *ctor2_o_fold2* =

trans[*OF* *conjunct2*[*OF* *fold_unique_mor*[*OF* *mor_comp*[*OF* *mor_fold* *mor_str*]]]] *fold2_ctor*

definition *dctor1* = *fold1* (*F1map* *id* *ctor1* *ctor2*) (*F2map* *id* *ctor1* *ctor2*)

definition *dctor2* = *fold2* (*F1map* *id* *ctor1* *ctor2*) (*F2map* *id* *ctor1* *ctor2*)

$\langle ML \rangle$

lemma *ctor1_o_dtor1*: $ctor1 \circ dtor1 = id$
 $\langle proof \rangle$

lemma *ctor2_o_dtor2*: $ctor2 \circ dtor2 = id$
 $\langle proof \rangle$

lemma *dtor1_o_ctor1*: $dtor1 \circ ctor1 = id$
 $\langle proof \rangle$

lemma *dtor2_o_ctor2*: $dtor2 \circ ctor2 = id$
 $\langle proof \rangle$

lemmas *dtor1_ctor1* = *pointfree_idE*[*OF dtor1_o_ctor1*]

lemmas *dtor2_ctor2* = *pointfree_idE*[*OF dtor2_o_ctor2*]

lemmas *ctor1_dtor1* = *pointfree_idE*[*OF ctor1_o_dtor1*]

lemmas *ctor2_dtor2* = *pointfree_idE*[*OF ctor2_o_dtor2*]

lemmas *bij_dtor1* = *o_bij*[*OF ctor1_o_dtor1 dtor1_o_ctor1*]

lemmas *inj_dtor1* = *bij_is_inj*[*OF bij_dtor1*]

lemmas *surj_dtor1* = *bij_is_surj*[*OF bij_dtor1*]

lemmas *dtor1_nchotomy* = *surjD*[*OF surj_dtor1*]

lemmas *dtor1_diff* = *inj_eq*[*OF inj_dtor1*]

lemmas *dtor1_cases* = *exE*[*OF dtor1_nchotomy*]

lemmas *bij_dtor2* = *o_bij*[*OF ctor2_o_dtor2 dtor2_o_ctor2*]

lemmas *inj_dtor2* = *bij_is_inj*[*OF bij_dtor2*]

lemmas *surj_dtor2* = *bij_is_surj*[*OF bij_dtor2*]

lemmas *dtor2_nchotomy* = *surjD*[*OF surj_dtor2*]

lemmas *dtor2_diff* = *inj_eq*[*OF inj_dtor2*]

lemmas *dtor2_cases* = *exE*[*OF dtor2_nchotomy*]

lemmas *bij_ctor1* = *o_bij*[*OF dtor1_o_ctor1 ctor1_o_dtor1*]

lemmas *inj_ctor1* = *bij_is_inj*[*OF bij_ctor1*]

lemmas *surj_ctor1* = *bij_is_surj*[*OF bij_ctor1*]

lemmas *ctor1_nchotomy* = *surjD*[*OF surj_ctor1*]

lemmas *ctor1_diff* = *inj_eq*[*OF inj_ctor1*]

lemmas *ctor1_cases* = *exE*[*OF ctor1_nchotomy*]

lemmas *bij_ctor2* = *o_bij*[*OF dtor2_o_ctor2 ctor2_o_dtor2*]

lemmas *inj_ctor2* = *bij_is_inj*[*OF bij_ctor2*]

lemmas *surj_ctor2* = *bij_is_surj*[*OF bij_ctor2*]

lemmas *ctor2_nchotomy* = *surjD*[*OF surj_ctor2*]

lemmas *ctor2_diff* = *inj_eq*[*OF inj_ctor2*]

lemmas *ctor2_cases* = *exE*[*OF ctor2_nchotomy*]

Primitive recursion

definition *rec1* **where**

$rec1\ s1\ s2 = snd\ o\ fold1\ (\langle ctor1\ o\ F1map\ id\ fst\ fst,\ s1 \rangle)\ (\langle ctor2\ o\ F2map\ id\ fst\ fst,\ s2 \rangle)$

definition *rec2* **where**

$rec2\ s1\ s2 = snd\ o\ fold2\ (\langle ctor1\ o\ F1map\ id\ fst\ fst,\ s1 \rangle)\ (\langle ctor2\ o\ F2map\ id\ fst\ fst,\ s2 \rangle)$

lemma *fold1_o_ctor1*: $fold1\ s1\ s2 \circ ctor1 = s1 \circ F1map\ id\ (fold1\ s1\ s2)\ (fold2\ s1\ s2)$
 $\langle proof \rangle$

lemma *fold2_o_ctor2*: $fold2\ s1\ s2 \circ ctor2 = s2 \circ F2map\ id\ (fold1\ s1\ s2)\ (fold2\ s1\ s2)$
 $\langle proof \rangle$

lemmas *fst_rec1_pair* =

$trans[OF\ conjunct1[OF\ fold_unique[OF$
 $\quad trans[OF\ o_assoc[symmetric]\ trans[OF\ arg_cong2[of\ _ _ _ _ (o),\ OF\ refl$
 $\quad \quad trans[OF\ fold1_o_ctor1\ convol_o]]],\ OF\ trans[OF\ fst_convol]]$
 $\quad trans[OF\ o_assoc[symmetric]\ trans[OF\ arg_cong2[of\ _ _ _ _ (o),\ OF\ refl$
 $\quad \quad trans[OF\ fold2_o_ctor2\ convol_o]]],\ OF\ trans[OF\ fst_convol]]]]$
 $fold1_ctor,\ unfolded\ F1.map_comp0[of\ id,\ unfolded\ id_o]\ F2.map_comp0[of\ id,\ unfolded\ id_o]\ o_assoc,$

$OF\ refl\ refl]$
lemmas $fst_rec2_pair =$
 $trans[OF\ conjunct2[OF\ fold_unique[OF$
 $\quad trans[OF\ o_assoc[symmetric]\ trans[OF\ arg_cong2[of\ _ _ _ _ (o),\ OF\ refl$
 $\quad\quad trans[OF\ fold1_o_ctor1\ convol_o]],\ OF\ trans[OF\ fst_convol]]$
 $\quad trans[OF\ o_assoc[symmetric]\ trans[OF\ arg_cong2[of\ _ _ _ _ (o),\ OF\ refl$
 $\quad\quad trans[OF\ fold2_o_ctor2\ convol_o]],\ OF\ trans[OF\ fst_convol]]]]]$
 $fold2_ctor,\ unfolded\ F1.map_comp0[of\ id,\ unfolded\ id_o]\ F2.map_comp0[of\ id,\ unfolded\ id_o]\ o_assoc,$
 $OF\ refl\ refl]$

theorem $rec1: rec1\ s1\ s2\ (ctor1\ x) = s1\ (F1map\ id\ (<id,\ rec1\ s1\ s2>)\ (<id,\ rec2\ s1\ s2>)\ x)$
 $\langle proof \rangle$

theorem $rec2: rec2\ s1\ s2\ (ctor2\ x) = s2\ (F2map\ id\ (<id,\ rec1\ s1\ s2>)\ (<id,\ rec2\ s1\ s2>)\ x)$
 $\langle proof \rangle$

lemma $rec_unique:$

$f\ o\ ctor1 = s1\ o\ F1map\ id\ <id,\ f>\ <id,\ g> \implies$
 $g\ o\ ctor2 = s2\ o\ F2map\ id\ <id,\ f>\ <id,\ g> \implies f = rec1\ s1\ s2 \wedge g = rec2\ s1\ s2$
 $\langle proof \rangle$

Induction

theorem $ctor_induct:$

$\llbracket \bigwedge x. (\bigwedge a. a \in F1set2\ x \implies phi1\ a) \implies (\bigwedge a. a \in F1set3\ x \implies phi2\ a) \implies phi1\ (ctor1\ x);$
 $\bigwedge x. (\bigwedge a. a \in F2set2\ x \implies phi1\ a) \implies (\bigwedge a. a \in F2set3\ x \implies phi2\ a) \implies phi2\ (ctor2\ x) \rrbracket \implies$
 $phi1\ a \wedge phi2\ b$
 $\langle proof \rangle$

theorem $ctor_induct2:$

$\llbracket \bigwedge x\ y. (\bigwedge a\ b. a \in F1set2\ x \implies b \in F1set2\ y \implies phi1\ a\ b) \implies$
 $(\bigwedge a\ b. a \in F1set3\ x \implies b \in F1set3\ y \implies phi2\ a\ b) \implies phi1\ (ctor1\ x)\ (ctor1\ y);$
 $\bigwedge x\ y. (\bigwedge a\ b. a \in F2set2\ x \implies b \in F2set2\ y \implies phi1\ a\ b) \implies$
 $(\bigwedge a\ b. a \in F2set3\ x \implies b \in F2set3\ y \implies phi2\ a\ b) \implies phi2\ (ctor2\ x)\ (ctor2\ y) \rrbracket \implies$
 $phi1\ a1\ b1 \wedge phi2\ a2\ b2$
 $\langle proof \rangle$

1.8 The Result as an BNF

The map operator

abbreviation $IF1map$ **where** $IF1map\ f \equiv fold1\ (ctor1\ o\ (F1map\ f\ id\ id))\ (ctor2\ o\ (F2map\ f\ id\ id))$

abbreviation $IF2map$ **where** $IF2map\ f \equiv fold2\ (ctor1\ o\ (F1map\ f\ id\ id))\ (ctor2\ o\ (F2map\ f\ id\ id))$

theorem $IF1map:$

$(IF1map\ f)\ o\ ctor1 = ctor1\ o\ (F1map\ f\ (IF1map\ f)\ (IF2map\ f))$
 $\langle proof \rangle$

theorem $IF2map:$

$(IF2map\ f)\ o\ ctor2 = ctor2\ o\ (F2map\ f\ (IF1map\ f)\ (IF2map\ f))$
 $\langle proof \rangle$

lemmas $IF1map_simps = o_eq_dest[OF\ IF1map]$

lemmas $IF2map_simps = o_eq_dest[OF\ IF2map]$

lemma $IFmap_unique:$

$\llbracket u\ o\ ctor1 = ctor1\ o\ F1map\ f\ u\ v;\ v\ o\ ctor2 = ctor2\ o\ F2map\ f\ u\ v \rrbracket \implies$
 $u = IF1map\ f \wedge v = IF2map\ f$
 $\langle proof \rangle$

theorem $IF1map_id: IF1map\ id = id$

$\langle proof \rangle$

theorem $IF2map_id: IF2map\ id = id$

$\langle proof \rangle$

theorem *IF1map_comp*: $IF1map (g \circ f) = IF1map g \circ IF1map f$
 ⟨proof⟩

theorem *IF2map_comp*: $IF2map (g \circ f) = IF2map g \circ IF2map f$
 ⟨proof⟩

The bound

abbreviation *IFbd* **where** $IFbd \equiv F1bd' +_c F2bd'$

theorem *IFbd_card_order*: $card_order\ IFbd$
 ⟨proof⟩

lemma *IFbd_Cinfinite*: $Cinfinite\ IFbd$
 ⟨proof⟩

lemmas *IFbd_cinfinite* = $conjunct1[OF\ IFbd_Cinfinite]$

The set operator

abbreviation *IF1col* **where** $IF1col \equiv (\lambda X. F1set1\ X \cup (\bigcup F1set2\ X \cup \bigcup F1set3\ X))$

abbreviation *IF2col* **where** $IF2col \equiv (\lambda X. F2set1\ X \cup (\bigcup F2set2\ X \cup \bigcup F2set3\ X))$

abbreviation *IF1set* **where** $IF1set \equiv fold1\ IF1col\ IF2col$

abbreviation *IF2set* **where** $IF2set \equiv fold2\ IF1col\ IF2col$

abbreviation *IF1in* **where** $IF1in\ A \equiv \{x. IF1set\ x \subseteq A\}$

abbreviation *IF2in* **where** $IF2in\ A \equiv \{x. IF2set\ x \subseteq A\}$

lemma *IF1set*: $IF1set\ o\ ctor1 = IF1col\ o\ (F1map\ id\ IF1set\ IF2set)$
 ⟨proof⟩

lemma *IF2set*: $IF2set\ o\ ctor2 = IF2col\ o\ (F2map\ id\ IF1set\ IF2set)$
 ⟨proof⟩

theorem *IF1set_simps*:
 $IF1set\ (ctor1\ x) = F1set1\ x \cup ((\bigcup a \in F1set2\ x. IF1set\ a) \cup (\bigcup a \in F1set3\ x. IF2set\ a))$
 ⟨proof⟩

theorem *IF2set_simps*:
 $IF2set\ (ctor2\ x) = F2set1\ x \cup ((\bigcup a \in F2set2\ x. IF1set\ a) \cup (\bigcup a \in F2set3\ x. IF2set\ a))$
 ⟨proof⟩

lemmas *F1set1_IF1set* = $xt1(3)[OF\ IF1set_simps\ Un_upper1]$

lemmas *F1set2_IF1set* = $subset_trans[OF\ UN_upper\ subset_trans[OF\ Un_upper1\ xt1(3)[OF\ IF1set_simps\ Un_upper2]]]$

lemmas *F1set3_IF1set* = $subset_trans[OF\ UN_upper\ subset_trans[OF\ Un_upper2\ xt1(3)[OF\ IF1set_simps\ Un_upper2]]]$

lemmas *F2set1_IF2set* = $xt1(3)[OF\ IF2set_simps\ Un_upper1]$

lemmas *F2set2_IF2set* = $subset_trans[OF\ UN_upper\ subset_trans[OF\ Un_upper1\ xt1(3)[OF\ IF2set_simps\ Un_upper2]]]$

lemmas *F2set3_IF2set* = $subset_trans[OF\ UN_upper\ subset_trans[OF\ Un_upper2\ xt1(3)[OF\ IF2set_simps\ Un_upper2]]]$

The BNF conditions for IF

lemma *IFset_natural*:
 $f' (IF1set\ x) = IF1set (IF1map\ f\ x) \wedge f' (IF2set\ y) = IF2set (IF2map\ f\ y)$
 ⟨proof⟩

theorem *IF1set_natural*: $IF1set\ o\ (IF1map\ f) = image\ f\ o\ IF1set$
 ⟨proof⟩

theorem *IF2set_natural*: $IF2set\ o\ (IF2map\ f) = image\ f\ o\ IF2set$
 ⟨proof⟩

lemma *IFmap_cong*:
 $(\forall a \in IF1set\ x. f\ a = g\ a) \longrightarrow IF1map\ f\ x = IF1map\ g\ x \wedge$

$(\forall a \in IF2set\ y. f\ a = g\ a) \longrightarrow IF2map\ f\ y = IF2map\ g\ y$
 ⟨proof⟩

theorem *IF1map_cong*:

$(\bigwedge a. a \in IF1set\ x \implies f\ a = g\ a) \implies IF1map\ f\ x = IF1map\ g\ x$
 ⟨proof⟩

theorem *IF2map_cong*:

$(\bigwedge a. a \in IF2set\ x \implies f\ a = g\ a) \implies IF2map\ f\ x = IF2map\ g\ x$
 ⟨proof⟩

lemma *IFset_bd*:

$|IF1set\ (x :: 'a\ IF1)| \leq_o\ IFbd \wedge |IF2set\ (y :: 'a\ IF2)| \leq_o\ IFbd$
 ⟨proof⟩

lemmas *IF1set_bd = conjunct1[OF IFset_bd]*

lemmas *IF2set_bd = conjunct2[OF IFset_bd]*

definition *IF1rel where*

$IF1rel\ R =$
 $(BNF_Def.Grp\ (IF1in\ (Collect\ (case_prod\ R)))\ (IF1map\ fst)) \wedge\ \dots \wedge$
 $(BNF_Def.Grp\ (IF1in\ (Collect\ (case_prod\ R)))\ (IF1map\ snd))$

definition *IF2rel where*

$IF2rel\ R =$
 $(BNF_Def.Grp\ (IF2in\ (Collect\ (case_prod\ R)))\ (IF2map\ fst)) \wedge\ \dots \wedge$
 $(BNF_Def.Grp\ (IF2in\ (Collect\ (case_prod\ R)))\ (IF2map\ snd))$

lemma *in_IF1rel*:

$IF1rel\ R\ x\ y \longleftrightarrow (\exists\ z. z \in IF1in\ (Collect\ (case_prod\ R)) \wedge IF1map\ fst\ z = x \wedge IF1map\ snd\ z = y)$
 ⟨proof⟩

lemma *in_IF2rel*:

$IF2rel\ R\ x\ y \longleftrightarrow (\exists\ z. z \in IF2in\ (Collect\ (case_prod\ R)) \wedge IF2map\ fst\ z = x \wedge IF2map\ snd\ z = y)$
 ⟨proof⟩

lemma *IF1rel_F1rel*: $IF1rel\ R\ (ctor1\ a)\ (ctor1\ b) \longleftrightarrow F1rel\ R\ (IF1rel\ R)\ (IF2rel\ R)\ a\ b$

⟨proof⟩

lemma *IF2rel_F2rel*: $IF2rel\ R\ (ctor2\ a)\ (ctor2\ b) \longleftrightarrow F2rel\ R\ (IF1rel\ R)\ (IF2rel\ R)\ a\ b$

⟨proof⟩

lemma *Irel_induct*:

assumes *IH1*: $\forall x\ y. F1rel\ P1\ P2\ P3\ x\ y \longrightarrow P2\ (ctor1\ x)\ (ctor1\ y)$

and *IH2*: $\forall x\ y. F2rel\ P1\ P2\ P3\ x\ y \longrightarrow P3\ (ctor2\ x)\ (ctor2\ y)$

shows $IF1rel\ P1 \leq P2 \wedge IF2rel\ P1 \leq P3$

⟨proof⟩

lemma *le_IFrel_Comp*:

$((IF1rel\ R\ OO\ IF1rel\ S)\ x1\ y1 \longrightarrow IF1rel\ (R\ OO\ S)\ x1\ y1) \wedge$
 $((IF2rel\ R\ OO\ IF2rel\ S)\ x2\ y2 \longrightarrow IF2rel\ (R\ OO\ S)\ x2\ y2)$

⟨proof⟩

lemma *le_IF1rel_Comp*: $IF1rel\ R1\ OO\ IF1rel\ R2 \leq IF1rel\ (R1\ OO\ R2)$

⟨proof⟩

lemma *le_IF2rel_Comp*: $IF2rel\ R1\ OO\ IF2rel\ R2 \leq IF2rel\ (R1\ OO\ R2)$

⟨proof⟩

context includes *lifting_syntax*

begin

lemma *fold_transfer*:

```

((F1rel R S T ==> S) ==> (F2rel R S T ==> T) ==> IF1rel R ==> S) fold1 fold1 ^
((F1rel R S T ==> S) ==> (F2rel R S T ==> T) ==> IF2rel R ==> T) fold2 fold2
⟨proof⟩

```

end

definition *IF1wit* $x = \text{ctor1 } (\text{wit2_F1 } x \text{ (ctor2 wit_F2)})$

definition *IF2wit* $= \text{ctor2 wit_F2}$

lemma *IF1wit*: $x \in \text{IF1set } (\text{IF1wit } y) \implies x = y$
 ⟨proof⟩

lemma *IF2wit*: $x \in \text{IF2set } \text{IF2wit} \implies \text{False}$
 ⟨proof⟩

⟨ML⟩

bnf 'a *IF1*
 map: *IF1map*
 sets: *IF1set*
 bd: *IFbd*
 wits: *IF1wit*
 rel: *IF1rel*
 ⟨proof⟩

bnf 'a *IF2*
 map: *IF2map*
 sets: *IF2set*
 bd: *IFbd*
 wits: *IF2wit*
 rel: *IF2rel*
 ⟨proof⟩

2 Greatest Fixpoint (a.k.a. Codatatype)

```

'b1 = ('a, 'b1, 'b2) F1
'b2 = ('a, 'b1, 'b2) F2

```

To build a witness scenario, let us assume

```

('a, 'b1, 'b2) F1 = 'a * 'b1 + 'a * 'b2
('a, 'b1, 'b2) F2 = unit + 'b1 * 'b2

```

⟨ML⟩

declare $[[\text{bnf_internals}]]$

bnf-axiomatization (*F1set1*: 'a, *F1set2*: 'b1, *F1set3*: 'b2) *F1*
 [wits: 'a \Rightarrow 'b1 \Rightarrow ('a, 'b1, 'b2) *F1* 'a \Rightarrow 'b2 \Rightarrow ('a, 'b1, 'b2) *F1*]
 for map: *F1map* rel: *F1rel*

bnf-axiomatization (*F2set1*: 'a, *F2set2*: 'b1, *F2set3*: 'b2) *F2*
 [wits: ('a, 'b1, 'b2) *F2*]
 for map: *F2map* rel: *F2rel*

lemma *F1rel_cong*: $[[R1 = S1; R2 = S2; R3 = S3]] \implies \text{F1rel } R1 \ R2 \ R3 = \text{F1rel } S1 \ S2 \ S3$
 ⟨proof⟩

lemma *F2rel_cong*: $[[R1 = S1; R2 = S2; R3 = S3]] \implies \text{F2rel } R1 \ R2 \ R3 = \text{F2rel } S1 \ S2 \ S3$
 ⟨proof⟩

abbreviation *F1in* :: 'a1 set \Rightarrow 'a2 set \Rightarrow 'a3 set \Rightarrow (('a1, 'a2, 'a3) *F1*) set **where**
F1in A1 A2 A3 $\equiv \{x. \text{F1set1 } x \subseteq A1 \wedge \text{F1set2 } x \subseteq A2 \wedge \text{F1set3 } x \subseteq A3\}$

abbreviation *F2in* :: 'a1 set \Rightarrow 'a2 set \Rightarrow 'a3 set \Rightarrow (('a1, 'a2, 'a3) *F2*) set **where**
F2in A1 A2 A3 $\equiv \{x. \text{F2set1 } x \subseteq A1 \wedge \text{F2set2 } x \subseteq A2 \wedge \text{F2set3 } x \subseteq A3\}$

lemma $F1map_comp_id$: $F1map\ g1\ g2\ g3\ (F1map\ id\ f2\ f3\ x) = F1map\ g1\ (g2\ o\ f2)\ (g3\ o\ f3)\ x$
 ⟨proof⟩

lemmas $F1in_mono23 = F1.in_mono[OF\ subset_refl]$

lemmas $F1in_mono23' = set_mp[OF\ F1in_mono23]$

lemma $F1map_congL$: $\llbracket \forall a \in F1set2\ x.\ f\ a = a; \forall a \in F1set3\ x.\ g\ a = a \rrbracket \implies$
 $F1map\ id\ f\ g\ x = x$
 ⟨proof⟩

lemma $F2map_comp_id$: $F2map\ g1\ g2\ g3\ (F2map\ id\ f2\ f3\ x) = F2map\ g1\ (g2\ o\ f2)\ (g3\ o\ f3)\ x$
 ⟨proof⟩

lemmas $F2in_mono23 = F2.in_mono[OF\ subset_refl]$

lemmas $F2in_mono23' = set_mp[OF\ F2in_mono23]$

lemma $F2map_congL$: $\llbracket \forall a \in F2set2\ x.\ f\ a = a; \forall a \in F2set3\ x.\ g\ a = a \rrbracket \implies$
 $F2map\ id\ f\ g\ x = x$
 ⟨proof⟩

2.1 Coalgebra

definition $coalg$ where

$coalg\ B1\ B2\ s1\ s2 =$
 $((\forall a \in B1.\ s1\ a \in F1in\ (UNIV :: 'a\ set)\ B1\ B2) \wedge (\forall a \in B2.\ s2\ a \in F2in\ (UNIV :: 'a\ set)\ B1\ B2))$

lemmas $coalg_F1in = bspec[OF\ conjunct1[OF\ iffD1[OF\ coalg_def]]]$

lemmas $coalg_F2in = bspec[OF\ conjunct2[OF\ iffD1[OF\ coalg_def]]]$

lemma $coalg_F1set2$:
 $\llbracket coalg\ B1\ B2\ s1\ s2; a \in B1 \rrbracket \implies F1set2\ (s1\ a) \subseteq B1$
 ⟨proof⟩

lemma $coalg_F1set3$:
 $\llbracket coalg\ B1\ B2\ s1\ s2; a \in B1 \rrbracket \implies F1set3\ (s1\ a) \subseteq B2$
 ⟨proof⟩

lemma $coalg_F2set2$:
 $\llbracket coalg\ B1\ B2\ s1\ s2; a \in B2 \rrbracket \implies F2set2\ (s2\ a) \subseteq B1$
 ⟨proof⟩

lemma $coalg_F2set3$:
 $\llbracket coalg\ B1\ B2\ s1\ s2; a \in B2 \rrbracket \implies F2set3\ (s2\ a) \subseteq B2$
 ⟨proof⟩

2.2 Type-coalgebra

abbreviation $tcoalg\ s1\ s2 \equiv coalg\ UNIV\ UNIV\ s1\ s2$

lemma $tcoalg$: $tcoalg\ s1\ s2$
 ⟨proof⟩

2.3 Morphism

definition mor where

$mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f\ g =$
 $((\forall a \in B1.\ f\ a \in B1') \wedge (\forall a \in B2.\ g\ a \in B2')) \wedge$
 $((\forall z \in B1.\ F1map\ (id :: 'a \Rightarrow 'a)\ f\ g\ (s1\ z) = s1'\ (f\ z)) \wedge$
 $(\forall z \in B2.\ F2map\ (id :: 'a \Rightarrow 'a)\ f\ g\ (s2\ z) = s2'\ (g\ z)))$

lemma mor_image1 : $mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f\ g \implies f\ ' B1 \subseteq B1'$
 ⟨proof⟩

lemma mor_image2 : $mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f\ g \implies g\ ' B2 \subseteq B2'$

<proof>

lemmas *mor_image1'* = *set_mp*[*OF mor_image1 imageI*]

lemmas *mor_image2'* = *set_mp*[*OF mor_image2 imageI*]

lemma *morE1*: $\llbracket \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g; \ z \in B1 \rrbracket$

$\implies F1map \ id \ f \ g \ (s1 \ z) = s1' \ (f \ z)$

<proof>

lemma *morE2*: $\llbracket \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g; \ z \in B2 \rrbracket$

$\implies F2map \ id \ f \ g \ (s2 \ z) = s2' \ (g \ z)$

<proof>

lemma *mor_incl*: $\llbracket B1 \subseteq B1'; \ B2 \subseteq B2' \rrbracket \implies \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1 \ s2 \ id \ id$

<proof>

lemmas *mor_id* = *mor_incl*[*OF subset_refl subset_refl*]

lemma *mor_comp*:

$\llbracket \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g;$

$\text{mor } B1' \ B2' \ s1' \ s2' \ B1'' \ B2'' \ s1'' \ s2'' \ f' \ g' \rrbracket \implies$

$\text{mor } B1 \ B2 \ s1 \ s2 \ B1'' \ B2'' \ s1'' \ s2'' \ (f' \circ f) \ (g' \circ g)$

<proof>

lemma *mor_cong*: $\llbracket f' = f; \ g' = g; \ \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g \rrbracket \implies$

$\text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f' \ g'$

<proof>

lemma *mor_UNIV*: $\text{mor } UNIV \ UNIV \ s1 \ s2 \ UNIV \ UNIV \ s1' \ s2' \ f1 \ f2 \longleftrightarrow$

$F1map \ id \ f1 \ f2 \circ s1 = s1' \circ f1 \wedge F2map \ id \ f1 \ f2 \circ s2 = s2' \circ f2$

<proof>

lemma *mor_str*:

$\text{mor } UNIV \ UNIV \ s1 \ s2 \ UNIV \ UNIV \ (F1map \ id \ s1 \ s2) \ (F2map \ id \ s1 \ s2) \ s1 \ s2$

<proof>

lemma *mor_case_sum*:

$\text{mor } UNIV \ UNIV \ s1 \ s2 \ UNIV \ UNIV \ (\text{case_sum } (F1map \ id \ Inl \ Inl \circ s1) \ s1') \ (\text{case_sum } (F2map \ id \ Inl \ Inl \circ s2) \ s2') \ Inl \ Inl$

<proof>

2.4 Bisimulations

definition *bis* where

$bis \ B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ R1 \ R2 =$

$((R1 \subseteq B1 \times B1' \wedge R2 \subseteq B2 \times B2') \wedge$

$(\forall b1 \ b1'. \ (b1, \ b1') \in R1 \longrightarrow$

$(\exists z \in F1in \ UNIV \ R1 \ R2.$

$F1map \ id \ fst \ fst \ z = s1 \ b1 \wedge F1map \ id \ snd \ snd \ z = s1' \ b1')) \wedge$

$(\forall b2 \ b2'. \ (b2, \ b2') \in R2 \longrightarrow$

$(\exists z \in F2in \ UNIV \ R1 \ R2.$

$F2map \ id \ fst \ fst \ z = s2 \ b2 \wedge F2map \ id \ snd \ snd \ z = s2' \ b2'))))$

lemma *bis_cong*: $\llbracket bis \ B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ R1 \ R2; \ R1' = R1; \ R2' = R2 \rrbracket \implies$

$bis \ B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ R1' \ R2'$

<proof>

lemma *bis_Frel*:

$bis \ B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ R1 \ R2 \longleftrightarrow$

$(R1 \subseteq B1 \times B1' \wedge R2 \subseteq B2 \times B2') \wedge$

$(\forall b1 \ b1'. \ (b1, \ b1') \in R1 \longrightarrow F1rel \ (=) \ (in_rel \ R1) \ (in_rel \ R2) \ (s1 \ b1) \ (s1' \ b1')) \wedge$

$(\forall b2 \ b2'. \ (b2, \ b2') \in R2 \longrightarrow F2rel \ (=) \ (in_rel \ R1) \ (in_rel \ R2) \ (s2 \ b2) \ (s2' \ b2'))$

<proof>

lemma *bis_converse*:

$bis\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ R1\ R2 \implies$
 $bis\ B1'\ B2'\ s1'\ s2'\ B1\ B2\ s1\ s2\ (R1 \hat{-} -1)\ (R2 \hat{-} -1)$
 ⟨proof⟩

lemma *bis_Comp*:

$\llbracket bis\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ P1\ P2;$
 $bis\ B1'\ B2'\ s1'\ s2'\ B1''\ B2''\ s1''\ s2''\ Q1\ Q2 \rrbracket \implies$
 $bis\ B1\ B2\ s1\ s2\ B1''\ B2''\ s1''\ s2''\ (P1\ O\ Q1)\ (P2\ O\ Q2)$
 ⟨proof⟩

lemma *bis_Gr*: $\llbracket coalg\ B1\ B2\ s1\ s2; mor\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f1\ f2 \rrbracket \implies$

$bis\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ (BNF_Def.Gr\ B1\ f1)\ (BNF_Def.Gr\ B2\ f2)$
 ⟨proof⟩

lemmas *bis_image2* = *bis_cong*[*OF bis_Comp*[*OF bis_converse*[*OF bis_Gr*] *bis_Gr*] *image2_Gr image2_Gr*]

lemmas *bis_diag* = *bis_cong*[*OF bis_Gr*[*OF _ mor_id*] *Id_on_Gr Id_on_Gr*]

lemma *bis_Union*: $\forall i \in I. bis\ B1\ B2\ s1\ s2\ B1\ B2\ s1\ s2\ (R1i\ i)\ (R2i\ i) \implies$

$bis\ B1\ B2\ s1\ s2\ B1\ B2\ s1\ s2\ (\bigcup_{i \in I} R1i\ i)\ (\bigcup_{i \in I} R2i\ i)$
 ⟨proof⟩

abbreviation *sbis* $B1\ B2\ s1\ s2\ R1\ R2 \equiv bis\ B1\ B2\ s1\ s2\ B1\ B2\ s1\ s2\ R1\ R2$

definition *lsbis1* **where** *lsbis1* $B1\ B2\ s1\ s2 =$

$(\bigcup R \in \{(R1, R2) \mid R1\ R2 . sbis\ B1\ B2\ s1\ s2\ R1\ R2\}. fst\ R)$

definition *lsbis2* **where** *lsbis2* $B1\ B2\ s1\ s2 =$

$(\bigcup R \in \{(R1, R2) \mid R1\ R2 . sbis\ B1\ B2\ s1\ s2\ R1\ R2\}. snd\ R)$

lemma *sbis_lsbis*:

$sbis\ B1\ B2\ s1\ s2\ (lsbis1\ B1\ B2\ s1\ s2)\ (lsbis2\ B1\ B2\ s1\ s2)$
 ⟨proof⟩

lemmas *lsbis1_incl* = *conjunct1*[*OF conjunct1*[*OF iffD1*[*OF bis_def*]], *OF sbis_lsbis*]

lemmas *lsbis2_incl* = *conjunct2*[*OF conjunct1*[*OF iffD1*[*OF bis_def*]], *OF sbis_lsbis*]

lemmas *lsbisE1* =

$mp[OF\ spec[OF\ spec[OF\ conjunct1[OF\ conjunct2[OF\ iffD1[OF\ bis_def]], OF\ sbis_lsbis]]]]$

lemmas *lsbisE2* =

$mp[OF\ spec[OF\ spec[OF\ conjunct2[OF\ conjunct2[OF\ iffD1[OF\ bis_def]], OF\ sbis_lsbis]]]]$

lemma *incl_lsbis1*: $sbis\ B1\ B2\ s1\ s2\ R1\ R2 \implies R1 \subseteq lsbis1\ B1\ B2\ s1\ s2$

⟨proof⟩

lemma *incl_lsbis2*: $sbis\ B1\ B2\ s1\ s2\ R1\ R2 \implies R2 \subseteq lsbis2\ B1\ B2\ s1\ s2$

⟨proof⟩

lemma *equiv_lsbis1*: $coalg\ B1\ B2\ s1\ s2 \implies equiv\ B1\ (lsbis1\ B1\ B2\ s1\ s2)$

⟨proof⟩

lemma *equiv_lsbis2*: $coalg\ B1\ B2\ s1\ s2 \implies equiv\ B2\ (lsbis2\ B1\ B2\ s1\ s2)$

⟨proof⟩

2.5 The Tree Coalgebra

typedef *bd_type_F* = *UNIV* :: $(bd_type_F1 + bd_type_F2)\ set$

⟨proof⟩

type-synonym *'a carrier* = $((bd_type_F + bd_type_F)\ list\ set \times$

$((bd_type_F + bd_type_F)\ list \Rightarrow ('a, bd_type_F, bd_type_F)\ F1 + ('a, bd_type_F, bd_type_F)\ F2))$

abbreviation *bd_F* $\equiv dir_image\ (bd_F1 + c\ bd_F2)\ Abs_bd_type_F$

lemmas $bd_F = \text{dir_image}[OF \text{Abs_bd_type_F_inject}[OF \text{UNIV_I UNIV_I}] \text{Card_order_csum}]$
lemmas $bd_F \text{Cinfinite} = \text{Cinfinite_cong}[OF \text{bd_F Cinfinite_csum1}[OF \text{F1.bd_Cinfinite}]]$
lemmas $bd_F \text{Card_order} = \text{Card_order_ordIso}[OF \text{Card_order_csum ordIso_symmetric}[OF \text{bd_F}]]$
lemma $bd_F \text{card_order}: \text{card_order } bd_F$
 (proof)

lemmas $F1set1_bd' = \text{ordLeq_transitive}[OF \text{F1.set_bd}(1) \text{ordLeq_ordIso_trans}[OF \text{ordLeq_csum1}[OF \text{F1.bd_Card_order}] \text{bd_F}]]$

lemmas $F1set2_bd' = \text{ordLeq_transitive}[OF \text{F1.set_bd}(2) \text{ordLeq_ordIso_trans}[OF \text{ordLeq_csum1}[OF \text{F1.bd_Card_order}] \text{bd_F}]]$

lemmas $F1set3_bd' = \text{ordLeq_transitive}[OF \text{F1.set_bd}(3) \text{ordLeq_ordIso_trans}[OF \text{ordLeq_csum1}[OF \text{F1.bd_Card_order}] \text{bd_F}]]$

lemmas $F2set1_bd' = \text{ordLeq_transitive}[OF \text{F2.set_bd}(1) \text{ordLeq_ordIso_trans}[OF \text{ordLeq_csum2}[OF \text{F2.bd_Card_order}] \text{bd_F}]]$

lemmas $F2set2_bd' = \text{ordLeq_transitive}[OF \text{F2.set_bd}(2) \text{ordLeq_ordIso_trans}[OF \text{ordLeq_csum2}[OF \text{F2.bd_Card_order}] \text{bd_F}]]$

lemmas $F2set3_bd' = \text{ordLeq_transitive}[OF \text{F2.set_bd}(3) \text{ordLeq_ordIso_trans}[OF \text{ordLeq_csum2}[OF \text{F2.bd_Card_order}] \text{bd_F}]]$

abbreviation $\text{Succ1 } Kl \text{ kl} \equiv \{k1. \text{Inl } k1 \in \text{BNF_Greatest_Fixpoint.Succ } Kl \text{ kl}\}$

abbreviation $\text{Succ2 } Kl \text{ kl} \equiv \{k2. \text{Inr } k2 \in \text{BNF_Greatest_Fixpoint.Succ } Kl \text{ kl}\}$

definition isNode1 where

$\text{isNode1 } Kl \text{ lab } kl = (\exists x1. \text{lab } kl = \text{Inl } x1 \wedge \text{F1set2 } x1 = \text{Succ1 } Kl \text{ kl} \wedge \text{F1set3 } x1 = \text{Succ2 } Kl \text{ kl})$

definition isNode2 where

$\text{isNode2 } Kl \text{ lab } kl = (\exists x2. \text{lab } kl = \text{Inr } x2 \wedge \text{F2set2 } x2 = \text{Succ1 } Kl \text{ kl} \wedge \text{F2set3 } x2 = \text{Succ2 } Kl \text{ kl})$

abbreviation isTree where

$\text{isTree } Kl \text{ lab} \equiv (\text{[]} \in Kl \wedge (\forall kl \in Kl. (\forall k1 \in \text{Succ1 } Kl \text{ kl}. \text{isNode1 } Kl \text{ lab } (kl \text{ @ } [\text{Inl } k1]))) \wedge (\forall k2 \in \text{Succ2 } Kl \text{ kl}. \text{isNode2 } Kl \text{ lab } (kl \text{ @ } [\text{Inr } k2])))$

definition carT1 where

$\text{carT1} = \{(Kl :: (\text{bd_type_F} + \text{bd_type_F}) \text{list set}, \text{lab}) \mid Kl \text{ lab}. \text{isTree } Kl \text{ lab} \wedge \text{isNode1 } Kl \text{ lab } \text{[]}\}$

definition carT2 where

$\text{carT2} = \{(Kl :: (\text{bd_type_F} + \text{bd_type_F}) \text{list set}, \text{lab}) \mid Kl \text{ lab}. \text{isTree } Kl \text{ lab} \wedge \text{isNode2 } Kl \text{ lab } \text{[]}\}$

definition strT1 where

$\text{strT1} = (\text{case_prod } (\%Kl \text{ lab}. \text{case_sum } (\text{F1map } id (\lambda k1. (\text{BNF_Greatest_Fixpoint.Shift } Kl (\text{Inl } k1), \text{BNF_Greatest_Fixpoint.shift } \text{lab } (\text{Inl } k1)))) (\lambda k2. (\text{BNF_Greatest_Fixpoint.Shift } Kl (\text{Inr } k2), \text{BNF_Greatest_Fixpoint.shift } \text{lab } (\text{Inr } k2)))) \text{undefined } (\text{lab } \text{[]}))$

definition strT2 where

$\text{strT2} = (\text{case_prod } (\%Kl \text{ lab}. \text{case_sum } \text{undefined } (\text{F2map } id (\lambda k1. (\text{BNF_Greatest_Fixpoint.Shift } Kl (\text{Inl } k1), \text{BNF_Greatest_Fixpoint.shift } \text{lab } (\text{Inl } k1)))) (\lambda k2. (\text{BNF_Greatest_Fixpoint.Shift } Kl (\text{Inr } k2), \text{BNF_Greatest_Fixpoint.shift } \text{lab } (\text{Inr } k2)))) (\text{lab } \text{[]}))$

lemma $\text{coalg_T}: \text{coalg } \text{carT1 } \text{carT2 } \text{strT1 } \text{strT2}$

(proof)

abbreviation tobd_F12 where $\text{tobd_F12 } s1 \text{ } x \equiv \text{toCard } (\text{F1set2 } (s1 \text{ } x)) \text{ } bd_F$

abbreviation tobd_F13 where $\text{tobd_F13 } s1 \text{ } x \equiv \text{toCard } (\text{F1set3 } (s1 \text{ } x)) \text{ } bd_F$

abbreviation tobd_F22 where $\text{tobd_F22 } s2 \text{ } x \equiv \text{toCard } (\text{F2set2 } (s2 \text{ } x)) \text{ } bd_F$

abbreviation tobd_F23 where $\text{tobd_F23 } s2 \text{ } x \equiv \text{toCard } (\text{F2set3 } (s2 \text{ } x)) \text{ } bd_F$

abbreviation frombd_F12 where $\text{frombd_F12 } s1 \text{ } x \equiv \text{fromCard } (\text{F1set2 } (s1 \text{ } x)) \text{ } bd_F$

abbreviation frombd_F13 where $\text{frombd_F13 } s1 \text{ } x \equiv \text{fromCard } (\text{F1set3 } (s1 \text{ } x)) \text{ } bd_F$

abbreviation frombd_F22 where $\text{frombd_F22 } s2 \text{ } x \equiv \text{fromCard } (\text{F2set2 } (s2 \text{ } x)) \text{ } bd_F$

abbreviation frombd_F23 where $\text{frombd_F23 } s2 \text{ } x \equiv \text{fromCard } (\text{F2set3 } (s2 \text{ } x)) \text{ } bd_F$

lemmas $\text{tobd_F12_inj} = \text{toCard_inj}[OF\ F1set2_bd'\ bd_F_Card_order]$
lemmas $\text{tobd_F13_inj} = \text{toCard_inj}[OF\ F1set3_bd'\ bd_F_Card_order]$
lemmas $\text{tobd_F22_inj} = \text{toCard_inj}[OF\ F2set2_bd'\ bd_F_Card_order]$
lemmas $\text{tobd_F23_inj} = \text{toCard_inj}[OF\ F2set3_bd'\ bd_F_Card_order]$
lemmas $\text{frombd_F12_tobd_F12} = \text{fromCard_toCard}[OF\ F1set2_bd'\ bd_F_Card_order]$
lemmas $\text{frombd_F13_tobd_F13} = \text{fromCard_toCard}[OF\ F1set3_bd'\ bd_F_Card_order]$
lemmas $\text{frombd_F22_tobd_F22} = \text{fromCard_toCard}[OF\ F2set2_bd'\ bd_F_Card_order]$
lemmas $\text{frombd_F23_tobd_F23} = \text{fromCard_toCard}[OF\ F2set3_bd'\ bd_F_Card_order]$

definition Lev where

$\text{Lev } s1\ s2 = \text{rec_nat } (\%a. \{\}, \%b. \{\})$
 $(\%n\ \text{rec.}$
 $(\%a1.$
 $\{ \text{Inl } (\text{tobd_F12 } s1\ a1\ b1) \# kl \mid b1\ kl. b1 \in F1set2\ (s1\ a1) \wedge kl \in \text{fst } \text{rec } b1 \} \cup$
 $\{ \text{Inr } (\text{tobd_F13 } s1\ a1\ b2) \# kl \mid b2\ kl. b2 \in F1set3\ (s1\ a1) \wedge kl \in \text{snd } \text{rec } b2 \},$
 $\%a2.$
 $\{ \text{Inl } (\text{tobd_F22 } s2\ a2\ b1) \# kl \mid b1\ kl. b1 \in F2set2\ (s2\ a2) \wedge kl \in \text{fst } \text{rec } b1 \} \cup$
 $\{ \text{Inr } (\text{tobd_F23 } s2\ a2\ b2) \# kl \mid b2\ kl. b2 \in F2set3\ (s2\ a2) \wedge kl \in \text{snd } \text{rec } b2 \})$

abbreviation Lev1 where $\text{Lev1 } s1\ s2\ n \equiv \text{fst } (\text{Lev } s1\ s2\ n)$

abbreviation Lev2 where $\text{Lev2 } s1\ s2\ n \equiv \text{snd } (\text{Lev } s1\ s2\ n)$

lemmas $\text{Lev1_0} = \text{fun_cong}[OF\ \text{fstI}[OF\ \text{rec_nat_0_imp}[OF\ \text{Lev_def}]]]$
lemmas $\text{Lev2_0} = \text{fun_cong}[OF\ \text{sndI}[OF\ \text{rec_nat_0_imp}[OF\ \text{Lev_def}]]]$
lemmas $\text{Lev1_Suc} = \text{fun_cong}[OF\ \text{fstI}[OF\ \text{rec_nat_Suc_imp}[OF\ \text{Lev_def}]]]$
lemmas $\text{Lev2_Suc} = \text{fun_cong}[OF\ \text{sndI}[OF\ \text{rec_nat_Suc_imp}[OF\ \text{Lev_def}]]]$

definition rv where

$\text{rv } s1\ s2 = \text{rec_list } (\%b1. \text{Inl } b1, \%b2. \text{Inr } b2)$
 $(\%k\ kl\ \text{rec.}$
 $(\%b1. \text{case_sum } (\%k1. \text{fst } \text{rec } (\text{frombd_F12 } s1\ b1\ k1)) (\%k2. \text{snd } \text{rec } (\text{frombd_F13 } s1\ b1\ k2))\ k,$
 $\%b2. \text{case_sum } (\%k1. \text{fst } \text{rec } (\text{frombd_F22 } s2\ b2\ k1)) (\%k2. \text{snd } \text{rec } (\text{frombd_F23 } s2\ b2\ k2))\ k))$

abbreviation rv1 where $\text{rv1 } s1\ s2\ kl \equiv \text{fst } (\text{rv } s1\ s2\ kl)$

abbreviation rv2 where $\text{rv2 } s1\ s2\ kl \equiv \text{snd } (\text{rv } s1\ s2\ kl)$

lemmas $\text{rv1_Nil} = \text{fun_cong}[OF\ \text{fstI}[OF\ \text{rec_list_Nil_imp}[OF\ \text{rv_def}]]]$
lemmas $\text{rv2_Nil} = \text{fun_cong}[OF\ \text{sndI}[OF\ \text{rec_list_Nil_imp}[OF\ \text{rv_def}]]]$
lemmas $\text{rv1_Cons} = \text{fun_cong}[OF\ \text{fstI}[OF\ \text{rec_list_Cons_imp}[OF\ \text{rv_def}]]]$
lemmas $\text{rv2_Cons} = \text{fun_cong}[OF\ \text{sndI}[OF\ \text{rec_list_Cons_imp}[OF\ \text{rv_def}]]]$

abbreviation $\text{Lab1 } s1\ s2\ b1\ kl \equiv$

$(\text{case_sum } (\%k. \text{Inl } (\text{F1map } \text{id } (\text{tobd_F12 } s1\ k) (\text{tobd_F13 } s1\ k) (s1\ k)))$
 $(\%k. \text{Inr } (\text{F2map } \text{id } (\text{tobd_F22 } s2\ k) (\text{tobd_F23 } s2\ k) (s2\ k))) (\text{rv1 } s1\ s2\ kl\ b1))$

abbreviation $\text{Lab2 } s1\ s2\ b2\ kl \equiv$

$(\text{case_sum } (\%k. \text{Inl } (\text{F1map } \text{id } (\text{tobd_F12 } s1\ k) (\text{tobd_F13 } s1\ k) (s1\ k)))$
 $(\%k. \text{Inr } (\text{F2map } \text{id } (\text{tobd_F22 } s2\ k) (\text{tobd_F23 } s2\ k) (s2\ k))) (\text{rv2 } s1\ s2\ kl\ b2))$

definition $\text{beh1 } s1\ s2\ a = (\bigcup n. \text{Lev1 } s1\ s2\ n\ a, \text{Lab1 } s1\ s2\ a)$

definition $\text{beh2 } s1\ s2\ a = (\bigcup n. \text{Lev2 } s1\ s2\ n\ a, \text{Lab2 } s1\ s2\ a)$

lemma length_Lev :

$\forall kl\ b1\ b2. (kl \in \text{Lev1 } s1\ s2\ n\ b1 \longrightarrow \text{length } kl = n) \wedge$
 $(kl \in \text{Lev2 } s1\ s2\ n\ b2 \longrightarrow \text{length } kl = n)$

$\langle \text{proof} \rangle$

lemmas $\text{length_Lev1} = \text{mp}[OF\ \text{conjunct1}[OF\ \text{spec}[OF\ \text{spec } [OF\ \text{spec}[OF\ \text{length_Lev}]]]]]$

lemmas $\text{length_Lev2} = \text{mp}[OF\ \text{conjunct2}[OF\ \text{spec}[OF\ \text{spec } [OF\ \text{spec}[OF\ \text{length_Lev}]]]]]$

lemma *length_Lev1'*: $kl \in Lev1\ s1\ s2\ n\ a \implies kl \in Lev1\ s1\ s2\ (length\ kl)\ a$
 ⟨proof⟩

lemma *length_Lev2'*: $kl \in Lev2\ s1\ s2\ n\ a \implies kl \in Lev2\ s1\ s2\ (length\ kl)\ a$
 ⟨proof⟩

lemma *rv_last*:

$\forall k\ b1\ b2.$
 $((\exists b1'.\ rv1\ s1\ s2\ (kl\ @\ [Inl\ k])\ b1 = Inl\ b1') \wedge$
 $(\exists b1'.\ rv1\ s1\ s2\ (kl\ @\ [Inr\ k])\ b1 = Inr\ b1')) \wedge$
 $((\exists b2'.\ rv2\ s1\ s2\ (kl\ @\ [Inl\ k])\ b2 = Inl\ b2') \wedge$
 $(\exists b2'.\ rv2\ s1\ s2\ (kl\ @\ [Inr\ k])\ b2 = Inr\ b2'))$
 ⟨proof⟩

lemmas *rv_last'* = *spec*[*OF spec*[*OF spec*[*OF spec*[*OF rv_last*]]]]

lemmas *rv1_Inl_last* = *conjunct1*[*OF conjunct1*[*OF rv_last'*]]

lemmas *rv1_Inr_last* = *conjunct2*[*OF conjunct1*[*OF rv_last'*]]

lemmas *rv2_Inl_last* = *conjunct1*[*OF conjunct2*[*OF rv_last'*]]

lemmas *rv2_Inr_last* = *conjunct2*[*OF conjunct2*[*OF rv_last'*]]

lemma *Fset_Lev*:

$\forall kl\ b1\ b2\ b1'\ b2'\ b1''\ b2''.$
 $(kl \in Lev1\ s1\ s2\ n\ b1 \implies$
 $((rv1\ s1\ s2\ kl\ b1 = Inl\ b1' \implies$
 $(b1'' \in F1set2\ (s1\ b1') \implies$
 $kl\ @\ [Inl\ (tobd_F12\ s1\ b1'\ b1'')] \in Lev1\ s1\ s2\ (Suc\ n)\ b1) \wedge$
 $(b2'' \in F1set3\ (s1\ b1') \implies$
 $kl\ @\ [Inr\ (tobd_F13\ s1\ b1'\ b2'')] \in Lev1\ s1\ s2\ (Suc\ n)\ b1)) \wedge$
 $(rv1\ s1\ s2\ kl\ b1 = Inr\ b2' \implies$
 $(b1'' \in F2set2\ (s2\ b2') \implies$
 $kl\ @\ [Inl\ (tobd_F22\ s2\ b2'\ b1'')] \in Lev1\ s1\ s2\ (Suc\ n)\ b1) \wedge$
 $(b2'' \in F2set3\ (s2\ b2') \implies$
 $kl\ @\ [Inr\ (tobd_F23\ s2\ b2'\ b2'')] \in Lev1\ s1\ s2\ (Suc\ n)\ b1)))) \wedge$
 $(kl \in Lev2\ s1\ s2\ n\ b2 \implies$
 $((rv2\ s1\ s2\ kl\ b2 = Inl\ b1' \implies$
 $(b1'' \in F1set2\ (s1\ b1') \implies$
 $kl\ @\ [Inl\ (tobd_F12\ s1\ b1'\ b1'')] \in Lev2\ s1\ s2\ (Suc\ n)\ b2) \wedge$
 $(b2'' \in F1set3\ (s1\ b1') \implies$
 $kl\ @\ [Inr\ (tobd_F13\ s1\ b1'\ b2'')] \in Lev2\ s1\ s2\ (Suc\ n)\ b2)) \wedge$
 $(rv2\ s1\ s2\ kl\ b2 = Inr\ b2' \implies$
 $(b1'' \in F2set2\ (s2\ b2') \implies$
 $kl\ @\ [Inl\ (tobd_F22\ s2\ b2'\ b1'')] \in Lev2\ s1\ s2\ (Suc\ n)\ b2) \wedge$
 $(b2'' \in F2set3\ (s2\ b2') \implies$
 $kl\ @\ [Inr\ (tobd_F23\ s2\ b2'\ b2'')] \in Lev2\ s1\ s2\ (Suc\ n)\ b2))))$
 ⟨proof⟩

lemmas *Fset_Lev'* = *spec*[*OF spec*[*OF spec*[*OF spec*[*OF spec*[*OF spec*[*OF spec*[*OF spec*[*OF Fset_Lev*]]]]]]]]]

lemmas *F1set2_Lev1* = *mp*[*OF conjunct1*[*OF mp*[*OF conjunct1*[*OF mp*[*OF conjunct1*[*OF Fset_Lev'*]]]]]]]]]

lemmas *F1set2_Lev2* = *mp*[*OF conjunct1*[*OF mp*[*OF conjunct1*[*OF mp*[*OF conjunct2*[*OF Fset_Lev'*]]]]]]]]]

lemmas *F2set2_Lev1* = *mp*[*OF conjunct1*[*OF mp*[*OF conjunct2*[*OF mp*[*OF conjunct1*[*OF Fset_Lev'*]]]]]]]]]

lemmas *F2set2_Lev2* = *mp*[*OF conjunct1*[*OF mp*[*OF conjunct2*[*OF mp*[*OF conjunct2*[*OF Fset_Lev'*]]]]]]]]]

lemmas *F1set3_Lev1* = *mp*[*OF conjunct2*[*OF mp*[*OF conjunct1*[*OF mp*[*OF conjunct1*[*OF Fset_Lev'*]]]]]]]]]

lemmas *F1set3_Lev2* = *mp*[*OF conjunct2*[*OF mp*[*OF conjunct1*[*OF mp*[*OF conjunct2*[*OF Fset_Lev'*]]]]]]]]]

lemmas *F2set3_Lev1* = *mp*[*OF conjunct2*[*OF mp*[*OF conjunct2*[*OF mp*[*OF conjunct1*[*OF Fset_Lev'*]]]]]]]]]

lemmas *F2set3_Lev2* = *mp*[*OF conjunct2*[*OF mp*[*OF conjunct2*[*OF mp*[*OF conjunct2*[*OF Fset_Lev'*]]]]]]]]]

lemma *Fset_image_Lev*:

$\forall kl\ k\ b1\ b2\ b1'\ b2'.$
 $(kl \in Lev1\ s1\ s2\ n\ b1 \implies$
 $(kl\ @\ [Inl\ k] \in Lev1\ s1\ s2\ (Suc\ n)\ b1 \implies$
 $(rv1\ s1\ s2\ kl\ b1 = Inl\ b1' \implies k \in tobd_F12\ s1\ b1'\ 'F1set2\ (s1\ b1')) \wedge$
 $(rv1\ s1\ s2\ kl\ b1 = Inr\ b2' \implies k \in tobd_F22\ s2\ b2'\ 'F2set2\ (s2\ b2')))) \wedge$

$(kl \ @ \ [Inr \ k] \in \text{Lev1 } s1 \ s2 \ (\text{Suc } n) \ b1 \ \longrightarrow$
 $(rv1 \ s1 \ s2 \ kl \ b1 = \text{Inl } b1' \ \longrightarrow \ k \in \text{tobd_F13 } s1 \ b1' \ ' \ \text{F1set3 } (s1 \ b1')) \wedge$
 $(rv1 \ s1 \ s2 \ kl \ b1 = \text{Inr } b2' \ \longrightarrow \ k \in \text{tobd_F23 } s2 \ b2' \ ' \ \text{F2set3 } (s2 \ b2')))) \wedge$
 $(kl \in \text{Lev2 } s1 \ s2 \ n \ b2 \ \longrightarrow$
 $(kl \ @ \ [\text{Inl } k] \in \text{Lev2 } s1 \ s2 \ (\text{Suc } n) \ b2 \ \longrightarrow$
 $(rv2 \ s1 \ s2 \ kl \ b2 = \text{Inl } b1' \ \longrightarrow \ k \in \text{tobd_F12 } s1 \ b1' \ ' \ \text{F1set2 } (s1 \ b1')) \wedge$
 $(rv2 \ s1 \ s2 \ kl \ b2 = \text{Inr } b2' \ \longrightarrow \ k \in \text{tobd_F22 } s2 \ b2' \ ' \ \text{F2set2 } (s2 \ b2')))) \wedge$
 $(kl \ @ \ [Inr \ k] \in \text{Lev2 } s1 \ s2 \ (\text{Suc } n) \ b2 \ \longrightarrow$
 $(rv2 \ s1 \ s2 \ kl \ b2 = \text{Inl } b1' \ \longrightarrow \ k \in \text{tobd_F13 } s1 \ b1' \ ' \ \text{F1set3 } (s1 \ b1')) \wedge$
 $(rv2 \ s1 \ s2 \ kl \ b2 = \text{Inr } b2' \ \longrightarrow \ k \in \text{tobd_F23 } s2 \ b2' \ ' \ \text{F2set3 } (s2 \ b2'))))$
 $\langle \text{proof} \rangle$

lemmas $\text{Fset_image_Lev}' =$
 $\text{spec}[OF \ \text{spec}[OF \ \text{spec}[OF \ \text{spec}[OF \ \text{spec}[OF \ \text{spec}[OF \ \text{Fset_image_Lev}]]]]]]]$
lemmas $\text{F1set2_image_Lev1} =$
 $\text{mp}[OF \ \text{conjunct1}[OF \ \text{mp}[OF \ \text{conjunct1}[OF \ \text{mp}[OF \ \text{conjunct1}[OF \ \text{Fset_image_Lev}]]]]]]]$
lemmas $\text{F1set2_image_Lev2} =$
 $\text{mp}[OF \ \text{conjunct1}[OF \ \text{mp}[OF \ \text{conjunct1}[OF \ \text{mp}[OF \ \text{conjunct2}[OF \ \text{Fset_image_Lev}]]]]]]]$
lemmas $\text{F1set3_image_Lev1} =$
 $\text{mp}[OF \ \text{conjunct1}[OF \ \text{mp}[OF \ \text{conjunct2}[OF \ \text{mp}[OF \ \text{conjunct1}[OF \ \text{Fset_image_Lev}]]]]]]]$
lemmas $\text{F1set3_image_Lev2} =$
 $\text{mp}[OF \ \text{conjunct1}[OF \ \text{mp}[OF \ \text{conjunct2}[OF \ \text{mp}[OF \ \text{conjunct2}[OF \ \text{Fset_image_Lev}]]]]]]]$
lemmas $\text{F2set2_image_Lev1} =$
 $\text{mp}[OF \ \text{conjunct2}[OF \ \text{mp}[OF \ \text{conjunct1}[OF \ \text{mp}[OF \ \text{conjunct1}[OF \ \text{Fset_image_Lev}]]]]]]]$
lemmas $\text{F2set2_image_Lev2} =$
 $\text{mp}[OF \ \text{conjunct2}[OF \ \text{mp}[OF \ \text{conjunct1}[OF \ \text{mp}[OF \ \text{conjunct2}[OF \ \text{Fset_image_Lev}]]]]]]]$
lemmas $\text{F2set3_image_Lev1} =$
 $\text{mp}[OF \ \text{conjunct2}[OF \ \text{mp}[OF \ \text{conjunct2}[OF \ \text{mp}[OF \ \text{conjunct1}[OF \ \text{Fset_image_Lev}]]]]]]]$
lemmas $\text{F2set3_image_Lev2} =$
 $\text{mp}[OF \ \text{conjunct2}[OF \ \text{mp}[OF \ \text{conjunct2}[OF \ \text{mp}[OF \ \text{conjunct2}[OF \ \text{Fset_image_Lev}]]]]]]]$

lemma $\text{mor_beh}:$
 $\text{mor } UNIV \ UNIV \ s1 \ s2 \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2} \ (\text{beh1 } \ s1 \ s2) \ (\text{beh2 } \ s1 \ s2)$
 $\langle \text{proof} \rangle$

2.6 Quotient Coalgebra

abbreviation car_final1 **where**
 $\text{car_final1} \equiv \text{carT1} \ // \ (\text{lsbis1 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2})$
abbreviation car_final2 **where**
 $\text{car_final2} \equiv \text{carT2} \ // \ (\text{lsbis2 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2})$
abbreviation str_final1 **where**
 $\text{str_final1} \equiv \text{univ } (\text{F1map } \ \text{id}$
 $(\text{Equiv_Relations.proj } (\text{lsbis1 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2}))$
 $(\text{Equiv_Relations.proj } (\text{lsbis2 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2})) \ o \ \text{strT1})$
abbreviation str_final2 **where**
 $\text{str_final2} \equiv \text{univ } (\text{F2map } \ \text{id}$
 $(\text{Equiv_Relations.proj } (\text{lsbis1 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2}))$
 $(\text{Equiv_Relations.proj } (\text{lsbis2 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2})) \ o \ \text{strT2})$

lemma $\text{congruent_strQ1}:$ $\text{congruent } (\text{lsbis1 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2} \ :: \ 'a \ \text{carrier } \ \text{rel})$
 $(\text{F1map } \ \text{id } (\text{Equiv_Relations.proj } (\text{lsbis1 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2} \ :: \ 'a \ \text{carrier } \ \text{rel})))$
 $(\text{Equiv_Relations.proj } (\text{lsbis2 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2} \ :: \ 'a \ \text{carrier } \ \text{rel})) \ o \ \text{strT1})$
 $\langle \text{proof} \rangle$

lemma $\text{congruent_strQ2}:$ $\text{congruent } (\text{lsbis2 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2} \ :: \ 'a \ \text{carrier } \ \text{rel})$
 $(\text{F2map } \ \text{id } (\text{Equiv_Relations.proj } (\text{lsbis1 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2} \ :: \ 'a \ \text{carrier } \ \text{rel})))$
 $(\text{Equiv_Relations.proj } (\text{lsbis2 } \ \text{carT1 } \ \text{carT2 } \ \text{strT1 } \ \text{strT2} \ :: \ 'a \ \text{carrier } \ \text{rel})) \ o \ \text{strT2})$
 $\langle \text{proof} \rangle$

lemma $\text{coalg_final}:$
 $\text{coalg } \text{car_final1 } \ \text{car_final2 } \ \text{str_final1 } \ \text{str_final2}$
 $\langle \text{proof} \rangle$

lemma *mor_T_final*:

mor carT1 carT2 strT1 strT2 car_final1 car_final2 str_final1 str_final2
(*Equiv_Relations.proj (lsbis1 carT1 carT2 strT1 strT2)*)
(*Equiv_Relations.proj (lsbis2 carT1 carT2 strT1 strT2)*)
{*proof*}

lemmas *mor_final = mor_comp*[*OF mor_beh mor_T_final*]

lemmas *in_car_final1 = mor_image1'*[*OF mor_final UNIV_I*]

lemmas *in_car_final2 = mor_image2'*[*OF mor_final UNIV_I*]

typedef (**overloaded**) 'a *JF1* = *car_final1* :: 'a carrier set set
{*proof*}

typedef (**overloaded**) 'a *JF2* = *car_final2* :: 'a carrier set set
{*proof*}

definition *dtor1* **where**

dtor1 x = F1map id Abs_JF1 Abs_JF2 (str_final1 (Rep_JF1 x))

definition *dtor2* **where**

dtor2 x = F2map id Abs_JF1 Abs_JF2 (str_final2 (Rep_JF2 x))

lemma *mor_Rep_JF: mor UNIV UNIV dtor1 dtor2*

car_final1 car_final2 str_final1 str_final2

Rep_JF1 Rep_JF2

{*proof*}

lemma *mor_Abs_JF: mor car_final1 car_final2 str_final1 str_final2*

UNIV UNIV dtor1 dtor2 Abs_JF1 Abs_JF2

{*proof*}

definition *unfold1* **where**

unfold1 s1 s2 x =

Abs_JF1 ((Equiv_Relations.proj (lsbis1 carT1 carT2 strT1 strT2) o beh1 s1 s2) x)

definition *unfold2* **where**

unfold2 s1 s2 x =

Abs_JF2 ((Equiv_Relations.proj (lsbis2 carT1 carT2 strT1 strT2) o beh2 s1 s2) x)

lemma *mor_unfold*:

mor UNIV UNIV s1 s2 UNIV UNIV dtor1 dtor2 (unfold1 s1 s2) (unfold2 s1 s2)

{*proof*}

lemmas *unfold1 = sym*[*OF morE1*[*OF mor_unfold UNIV_I*]]

lemmas *unfold2 = sym*[*OF morE2*[*OF mor_unfold UNIV_I*]]

lemma *JF_cind: sbis UNIV UNIV dtor1 dtor2 R1 R2 $\implies R1 \subseteq Id \wedge R2 \subseteq Id$*

{*proof*}

lemmas *JF_cind1 = conjunct1*[*OF JF_cind*]

lemmas *JF_cind2 = conjunct2*[*OF JF_cind*]

lemma *unfold_unique_mor*:

mor UNIV UNIV s1 s2 UNIV UNIV dtor1 dtor2 f1 f2 \implies

f1 = unfold1 s1 s2 \wedge f2 = unfold2 s1 s2

{*proof*}

lemmas *unfold_unique = unfold_unique_mor*[*OF iffD2*[*OF mor_UNIV*], *OF conjI*]

lemmas *unfold1_dtor = sym*[*OF conjunct1*[*OF unfold_unique_mor*[*OF mor_id*]]]

lemmas *unfold2_dtor = sym*[*OF conjunct2*[*OF unfold_unique_mor*[*OF mor_id*]]]

lemmas *unfold1_o_dtor1* =
trans[*OF conjunct1*[*OF unfold_unique_mor*[*OF mor_comp*[*OF mor_str mor_unfold*]]] *unfold1_dtor*]
lemmas *unfold2_o_dtor2* =
trans[*OF conjunct2*[*OF unfold_unique_mor*[*OF mor_comp*[*OF mor_str mor_unfold*]]] *unfold2_dtor*]

definition *ctor1* **where** *ctor1* = *unfold1* (*F1map id dtor1 dtor2*) (*F2map id dtor1 dtor2*)
definition *ctor2* **where** *ctor2* = *unfold2* (*F1map id dtor1 dtor2*) (*F2map id dtor1 dtor2*)

lemma *ctor1_o_dtor1*:
ctor1 o dtor1 = id
<proof>

lemma *ctor2_o_dtor2*:
ctor2 o dtor2 = id
<proof>

lemma *dtor1_o_ctor1*:
dtor1 o ctor1 = id
<proof>

lemma *dtor2_o_ctor2*:
dtor2 o ctor2 = id
<proof>

lemmas *dtor1_ctor1 = pointfree_idE*[*OF dtor1_o_ctor1*]
lemmas *dtor2_ctor2 = pointfree_idE*[*OF dtor2_o_ctor2*]
lemmas *ctor1_dtor1 = pointfree_idE*[*OF ctor1_o_dtor1*]
lemmas *ctor2_dtor2 = pointfree_idE*[*OF ctor2_o_dtor2*]

lemmas *bij_dtor1 = o_bij*[*OF ctor1_o_dtor1 dtor1_o_ctor1*]
lemmas *inj_dtor1 = bij_is_inj*[*OF bij_dtor1*]
lemmas *surj_dtor1 = bij_is_surj*[*OF bij_dtor1*]
lemmas *dtor1_nchotomy = surjD*[*OF surj_dtor1*]
lemmas *dtor1_diff = inj_eq*[*OF inj_dtor1*]
lemmas *dtor1_cases = exE*[*OF dtor1_nchotomy*]
lemmas *bij_dtor2 = o_bij*[*OF ctor2_o_dtor2 dtor2_o_ctor2*]
lemmas *inj_dtor2 = bij_is_inj*[*OF bij_dtor2*]
lemmas *surj_dtor2 = bij_is_surj*[*OF bij_dtor2*]
lemmas *dtor2_nchotomy = surjD*[*OF surj_dtor2*]
lemmas *dtor2_diff = inj_eq*[*OF inj_dtor2*]
lemmas *dtor2_cases = exE*[*OF dtor2_nchotomy*]

lemmas *bij_ctor1 = o_bij*[*OF dtor1_o_ctor1 ctor1_o_dtor1*]
lemmas *inj_ctor1 = bij_is_inj*[*OF bij_ctor1*]
lemmas *surj_ctor1 = bij_is_surj*[*OF bij_ctor1*]
lemmas *ctor1_nchotomy = surjD*[*OF surj_ctor1*]
lemmas *ctor1_diff = inj_eq*[*OF inj_ctor1*]
lemmas *ctor1_cases = exE*[*OF ctor1_nchotomy*]
lemmas *bij_ctor2 = o_bij*[*OF dtor2_o_ctor2 ctor2_o_dtor2*]
lemmas *inj_ctor2 = bij_is_inj*[*OF bij_ctor2*]
lemmas *surj_ctor2 = bij_is_surj*[*OF bij_ctor2*]
lemmas *ctor2_nchotomy = surjD*[*OF surj_ctor2*]
lemmas *ctor2_diff = inj_eq*[*OF inj_ctor2*]
lemmas *ctor2_cases = exE*[*OF ctor2_nchotomy*]

lemmas *ctor1_unfold1 = iffD1*[*OF dtor1_diff trans*[*OF unfold1 sym*[*OF dtor1_ctor1*]]]
lemmas *ctor2_unfold2 = iffD1*[*OF dtor2_diff trans*[*OF unfold2 sym*[*OF dtor2_ctor2*]]]

definition *corec1* **where** *corec1 s1 s2 =*
unfold1 (*case_sum* (*F1map id Inl Inl o dtor1*) *s1*)

$(\text{case_sum } (F2\text{map } \text{id } \text{Inl } \text{Inl } \circ \text{dtor2}) \text{ } s2) \circ \text{Inr}$
definition *corec2* **where** *corec2* $s1 \text{ } s2 =$
 $\text{unfold2 } (\text{case_sum } (F1\text{map } \text{id } \text{Inl } \text{Inl } \circ \text{dtor1}) \text{ } s1)$
 $(\text{case_sum } (F2\text{map } \text{id } \text{Inl } \text{Inl } \circ \text{dtor2}) \text{ } s2) \circ \text{Inr}$

lemma *dtor1_o_unfold1*: $\text{dtor1 } \circ \text{unfold1 } s1 \text{ } s2 = F1\text{map } \text{id } (\text{unfold1 } s1 \text{ } s2) (\text{unfold2 } s1 \text{ } s2) \circ s1$
 $\langle \text{proof} \rangle$

lemma *dtor2_o_unfold2*: $\text{dtor2 } \circ \text{unfold2 } s1 \text{ } s2 = F2\text{map } \text{id } (\text{unfold1 } s1 \text{ } s2) (\text{unfold2 } s1 \text{ } s2) \circ s2$
 $\langle \text{proof} \rangle$

lemma *corec1_Inl_sum*:
 $\text{unfold1 } (\text{case_sum } (F1\text{map } \text{id } \text{Inl } \text{Inl } \circ \text{dtor1}) \text{ } s1) (\text{case_sum } (F2\text{map } \text{id } \text{Inl } \text{Inl } \circ \text{dtor2}) \text{ } s2) \circ \text{Inl} = \text{id}$
 $\langle \text{proof} \rangle$

lemma *corec2_Inl_sum*:
 $\text{unfold2 } (\text{case_sum } (F1\text{map } \text{id } \text{Inl } \text{Inl } \circ \text{dtor1}) \text{ } s1) (\text{case_sum } (F2\text{map } \text{id } \text{Inl } \text{Inl } \circ \text{dtor2}) \text{ } s2) \circ \text{Inl} = \text{id}$
 $\langle \text{proof} \rangle$

lemma *case_sum_expand_Inr*: $f \circ \text{Inl} = g \implies \text{case_sum } g (f \circ \text{Inr}) = f$
 $\langle \text{proof} \rangle$

theorem *corec1*:
 $\text{dtor1 } (\text{corec1 } s1 \text{ } s2 \text{ } a) =$
 $F1\text{map } \text{id } (\text{case_sum } \text{id } (\text{corec1 } s1 \text{ } s2)) (\text{case_sum } \text{id } (\text{corec2 } s1 \text{ } s2)) (s1 \text{ } a)$
 $\langle \text{proof} \rangle$

theorem *corec2*:
 $\text{dtor2 } (\text{corec2 } s1 \text{ } s2 \text{ } a) =$
 $F2\text{map } \text{id } (\text{case_sum } \text{id } (\text{corec1 } s1 \text{ } s2)) (\text{case_sum } \text{id } (\text{corec2 } s1 \text{ } s2)) (s2 \text{ } a)$
 $\langle \text{proof} \rangle$

lemma *corec_unique*:
 $F1\text{map } \text{id } (\text{case_sum } \text{id } f1) (\text{case_sum } \text{id } f2) \circ s1 = \text{dtor1 } \circ f1 \implies$
 $F2\text{map } \text{id } (\text{case_sum } \text{id } f1) (\text{case_sum } \text{id } f2) \circ s2 = \text{dtor2 } \circ f2 \implies$
 $f1 = \text{corec1 } s1 \text{ } s2 \wedge f2 = \text{corec2 } s1 \text{ } s2$
 $\langle \text{proof} \rangle$

2.7 Coinduction

lemma *Frel_coind*:
 $\llbracket \forall a \text{ } b. \text{phi1 } a \text{ } b \longrightarrow F1\text{rel } (=) \text{phi1 } \text{phi2 } (\text{dtor1 } a) (\text{dtor1 } b);$
 $\forall a \text{ } b. \text{phi2 } a \text{ } b \longrightarrow F2\text{rel } (=) \text{phi1 } \text{phi2 } (\text{dtor2 } a) (\text{dtor2 } b) \rrbracket \implies$
 $(\text{phi1 } a1 \text{ } b1 \longrightarrow a1 = b1) \wedge (\text{phi2 } a2 \text{ } b2 \longrightarrow a2 = b2)$
 $\langle \text{proof} \rangle$

2.8 The Result as an BNF

abbreviation *JF1map* **where**
 $JF1\text{map } u \equiv \text{unfold1 } (F1\text{map } u \text{ } \text{id } \text{id } \circ \text{dtor1}) (F2\text{map } u \text{ } \text{id } \text{id } \circ \text{dtor2})$

abbreviation *JF2map* **where**
 $JF2\text{map } u \equiv \text{unfold2 } (F1\text{map } u \text{ } \text{id } \text{id } \circ \text{dtor1}) (F2\text{map } u \text{ } \text{id } \text{id } \circ \text{dtor2})$

lemma *JF1map*: $\text{dtor1 } \circ JF1\text{map } u = F1\text{map } u (JF1\text{map } u) (JF2\text{map } u) \circ \text{dtor1}$
 $\langle \text{proof} \rangle$

lemma *JF2map*: $\text{dtor2 } \circ JF2\text{map } u = F2\text{map } u (JF1\text{map } u) (JF2\text{map } u) \circ \text{dtor2}$
 $\langle \text{proof} \rangle$

lemmas *JF1map_simps* = $\text{o_eq_dest}[OF JF1\text{map}]$

lemmas *JF2map_simps* = $\text{o_eq_dest}[OF JF2\text{map}]$

theorem *JF1map_id*: $JF1\text{map } \text{id} = \text{id}$
 $\langle \text{proof} \rangle$

theorem *JF2map_id*: $JF2map\ id = id$
 ⟨proof⟩

lemma *JFmap_unique*:
 $\llbracket dtor1\ o\ u = F1map\ f\ u\ v\ o\ dtor1; dtor2\ o\ v = F2map\ f\ u\ v\ o\ dtor2 \rrbracket \implies$
 $u = JF1map\ f \wedge v = JF2map\ f$
 ⟨proof⟩

theorem *JF1map_comp*: $JF1map\ (g\ o\ f) = JF1map\ g\ o\ JF1map\ f$
 ⟨proof⟩

theorem *JF2map_comp*: $JF2map\ (g\ o\ f) = JF2map\ g\ o\ JF2map\ f$
 ⟨proof⟩

definition *JFcol* **where**

$JFcol = rec_nat\ (\%a.\ \{\},\ \%b.\ \{\})$
 (%n rec.
 (%a. $F1set1\ (dtor1\ a) \cup$
 $((\bigcup a' \in F1set2\ (dtor1\ a).\ fst\ rec\ a') \cup$
 $(\bigcup a' \in F1set3\ (dtor1\ a).\ snd\ rec\ a'))$),
 %b. $F2set1\ (dtor2\ b) \cup$
 $((\bigcup b' \in F2set2\ (dtor2\ b).\ fst\ rec\ b') \cup$
 $(\bigcup b' \in F2set3\ (dtor2\ b).\ snd\ rec\ b'))$))

abbreviation *JF1col* **where** $JF1col\ n \equiv fst\ (JFcol\ n)$

abbreviation *JF2col* **where** $JF2col\ n \equiv snd\ (JFcol\ n)$

lemmas *JF1col_0* = $fun_cong[OF\ fstI[OF\ rec_nat_0_imp[OF\ JFcol_def]]]$

lemmas *JF2col_0* = $fun_cong[OF\ sndI[OF\ rec_nat_0_imp[OF\ JFcol_def]]]$

lemmas *JF1col_Suc* = $fun_cong[OF\ fstI[OF\ rec_nat_Suc_imp[OF\ JFcol_def]]]$

lemmas *JF2col_Suc* = $fun_cong[OF\ sndI[OF\ rec_nat_Suc_imp[OF\ JFcol_def]]]$

lemma *JFcol_minimal*:

$\llbracket \bigwedge a.\ F1set1\ (dtor1\ a) \subseteq K1\ a;$
 $\bigwedge b.\ F2set1\ (dtor2\ b) \subseteq K2\ b;$
 $\bigwedge a\ a'.\ a' \in F1set2\ (dtor1\ a) \implies K1\ a' \subseteq K1\ a;$
 $\bigwedge a\ b'.\ b' \in F1set3\ (dtor1\ a) \implies K2\ b' \subseteq K1\ a;$
 $\bigwedge b\ a'.\ a' \in F2set2\ (dtor2\ b) \implies K1\ a' \subseteq K2\ b;$
 $\bigwedge b\ b'.\ b' \in F2set3\ (dtor2\ b) \implies K2\ b' \subseteq K2\ b \rrbracket \implies$
 $\forall a\ b.\ JF1col\ n\ a \subseteq K1\ a \wedge JF2col\ n\ b \subseteq K2\ b$
 ⟨proof⟩

lemma *JFset_minimal*:

$\llbracket \bigwedge a.\ F1set1\ (dtor1\ a) \subseteq K1\ a;$
 $\bigwedge b.\ F2set1\ (dtor2\ b) \subseteq K2\ b;$
 $\bigwedge a\ a'.\ a' \in F1set2\ (dtor1\ a) \implies K1\ a' \subseteq K1\ a;$
 $\bigwedge a\ b'.\ b' \in F1set3\ (dtor1\ a) \implies K2\ b' \subseteq K1\ a;$
 $\bigwedge b\ a'.\ a' \in F2set2\ (dtor2\ b) \implies K1\ a' \subseteq K2\ b;$
 $\bigwedge b\ b'.\ b' \in F2set3\ (dtor2\ b) \implies K2\ b' \subseteq K2\ b \rrbracket \implies$
 $(\bigcup n.\ JF1col\ n\ a) \subseteq K1\ a \wedge (\bigcup n.\ JF2col\ n\ b) \subseteq K2\ b$
 ⟨proof⟩

abbreviation *JF1set* **where** $JF1set\ a \equiv (\bigcup n.\ JF1col\ n\ a)$

abbreviation *JF2set* **where** $JF2set\ a \equiv (\bigcup n.\ JF2col\ n\ a)$

lemma *F1set1_incl_JF1set*:

$F1set1\ (dtor1\ a) \subseteq JF1set\ a$
 ⟨proof⟩

lemma *F2set1_incl_JF2set*:

$F2set1\ (dtor2\ a) \subseteq JF2set\ a$
 ⟨proof⟩

lemma *F1set2_JF1set_incl_JF1set*:
 $a' \in F1set2 \text{ (dtor1 } a) \implies JF1set \ a' \subseteq JF1set \ a$
 ⟨proof⟩

lemma *F1set3_JF2set_incl_JF1set*:
 $a' \in F1set3 \text{ (dtor1 } a) \implies JF2set \ a' \subseteq JF1set \ a$
 ⟨proof⟩

lemma *F2set2_JF1set_incl_JF2set*:
 $a' \in F2set2 \text{ (dtor2 } a) \implies JF1set \ a' \subseteq JF2set \ a$
 ⟨proof⟩

lemma *F2set3_JF2set_incl_JF2set*:
 $a' \in F2set3 \text{ (dtor2 } a) \implies JF2set \ a' \subseteq JF2set \ a$
 ⟨proof⟩

lemmas *F1set1_JF1set = set_mp[OF F1set1_incl_JF1set]*
lemmas *F2set1_JF2set = set_mp[OF F2set1_incl_JF2set]*
lemmas *F1set2_JF1set_JF1set = set_mp[OF F1set2_JF1set_incl_JF1set]*
lemmas *F1set3_JF2set_JF1set = set_mp[OF F1set3_JF2set_incl_JF1set]*
lemmas *F2set2_JF1set_JF2set = set_mp[OF F2set2_JF1set_incl_JF2set]*
lemmas *F2set3_JF2set_JF2set = set_mp[OF F2set3_JF2set_incl_JF2set]*

lemma *JFset_le*:
fixes $a :: 'a \ JF1$ **and** $b :: 'a \ JF2$
shows
 $JF1set \ a \subseteq F1set1 \text{ (dtor1 } a) \cup (UNION \ (F1set2 \text{ (dtor1 } a)) \ JF1set \cup UNION \ (F1set3 \text{ (dtor1 } a)) \ JF2set) \wedge$
 $JF2set \ b \subseteq F2set1 \text{ (dtor2 } b) \cup (UNION \ (F2set2 \text{ (dtor2 } b)) \ JF1set \cup UNION \ (F2set3 \text{ (dtor2 } b)) \ JF2set)$
 ⟨proof⟩

theorem *JF1set_simps*:
 $JF1set \ a = F1set1 \text{ (dtor1 } a) \cup$
 $((\bigcup \ b \in F1set2 \text{ (dtor1 } a). \ JF1set \ b) \cup$
 $(\bigcup \ b \in F1set3 \text{ (dtor1 } a). \ JF2set \ b))$
 ⟨proof⟩

theorem *JF2set_simps*:
 $JF2set \ a = F2set1 \text{ (dtor2 } a) \cup$
 $((\bigcup \ b \in F2set2 \text{ (dtor2 } a). \ JF1set \ b) \cup$
 $(\bigcup \ b \in F2set3 \text{ (dtor2 } a). \ JF2set \ b))$
 ⟨proof⟩

lemma *JFcol_natural*:
 $\forall b1 \ b2. \ u \ ' \ (JF1col \ n \ b1) = JF1col \ n \ (JF1map \ u \ b1) \wedge$
 $u \ ' \ (JF2col \ n \ b2) = JF2col \ n \ (JF2map \ u \ b2)$
 ⟨proof⟩

theorem *JF1set_natural*: $JF1set \ o \ (JF1map \ u) = image \ u \ o \ JF1set$
 ⟨proof⟩

theorem *JF2set_natural*: $JF2set \ o \ (JF2map \ u) = image \ u \ o \ JF2set$
 ⟨proof⟩

theorem *JFmap_cong0*:
 $((\forall p \in JF1set \ a. \ u \ p = v \ p) \longrightarrow JF1map \ u \ a = JF1map \ v \ a) \wedge$
 $((\forall p \in JF2set \ b. \ u \ p = v \ p) \longrightarrow JF2map \ u \ b = JF2map \ v \ b)$
 ⟨proof⟩

lemmas *JF1map_cong0 = mp[OF conjunct1[OF JFmap_cong0]]*
lemmas *JF2map_cong0 = mp[OF conjunct2[OF JFmap_cong0]]*

lemma *JFcol_bd*: $\forall (j1 :: 'a \ JF1) \ (j2 :: 'a \ JF2). \ |JF1col \ n \ j1| \leq o \ bd_F \wedge \ |JF2col \ n \ j2| \leq o \ bd_F$

<proof>

theorem *JF1set_bd*: $|JF1set\ j| \leq o\ bd_F$
<proof>

theorem *JF2set_bd*: $|JF2set\ j| \leq o\ bd_F$
<proof>

abbreviation *JF2wit* \equiv *ctor2 wit_F2*

theorem *JF2wit*: $\bigwedge x. x \in JF2set\ JF2wit \implies False$
<proof>

abbreviation *JF1wit* \equiv ($\%a. ctor1\ (wit2_F1\ a\ JF2wit)$)

theorem *JF1wit*: $\bigwedge x. x \in JF1set\ (JF1wit\ a) \implies x = a$
<proof>

context

fixes *phi1* :: $'a \Rightarrow 'a\ JF1 \Rightarrow bool$ **and** *phi2* :: $'a \Rightarrow 'a\ JF2 \Rightarrow bool$

begin

lemmas *JFset_induct* =

JFset_minimal[of $\%b1. \{a \in JF1set\ b1 . phi1\ a\ b1\} \%b2. \{a \in JF2set\ b2 . phi2\ a\ b2\}$,
unfolded subset_Collect_iff[*OF F1set1_incl_JF1set*] *subset_Collect_iff*[*OF F2set1_incl_JF2set*]
subset_Collect_iff[*OF subset_refl*],
OF ballI ballI
subset_CollectI[*OF F1set2_JF1set_incl_JF1set*]
subset_CollectI[*OF F1set3_JF2set_incl_JF1set*]
subset_CollectI[*OF F2set2_JF1set_incl_JF2set*]
subset_CollectI[*OF F2set3_JF2set_incl_JF2set*]]

end

<ML>

abbreviation *JF1in* **where** *JF1in* $B \equiv \{a. JF1set\ a \subseteq B\}$

abbreviation *JF2in* **where** *JF2in* $B \equiv \{a. JF2set\ a \subseteq B\}$

definition *JF1rel* **where**

JF1rel $R = (BNF_Def.Grp\ (JF1in\ (Collect\ (case_prod\ R)))\ (JF1map\ fst))^{--1}\ OO$
 $(BNF_Def.Grp\ (JF1in\ (Collect\ (case_prod\ R)))\ (JF1map\ snd))$

definition *JF2rel* **where**

JF2rel $R = (BNF_Def.Grp\ (JF2in\ (Collect\ (case_prod\ R)))\ (JF2map\ fst))^{--1}\ OO$
 $(BNF_Def.Grp\ (JF2in\ (Collect\ (case_prod\ R)))\ (JF2map\ snd))$

lemma *in_JF1rel*:

JF1rel $R\ x\ y \longleftrightarrow (\exists z. z \in JF1in\ (Collect\ (case_prod\ R)) \wedge JF1map\ fst\ z = x \wedge JF1map\ snd\ z = y)$
<proof>

lemma *in_JF2rel*:

JF2rel $R\ x\ y \longleftrightarrow (\exists z. z \in JF2in\ (Collect\ (case_prod\ R)) \wedge JF2map\ fst\ z = x \wedge JF2map\ snd\ z = y)$
<proof>

lemma *J_rel_coind_ind*:

$\llbracket \forall x y. R2\ x\ y \longrightarrow (f\ x\ y \in F1in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))) \wedge$
 $F1map\ fst\ fst\ fst\ (f\ x\ y) = dtor1\ x \wedge$
 $F1map\ snd\ snd\ snd\ (f\ x\ y) = dtor1\ y;$
 $\forall x y. R3\ x\ y \longrightarrow (g\ x\ y \in F2in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))) \wedge$
 $F2map\ fst\ fst\ fst\ (g\ x\ y) = dtor2\ x \wedge$
 $F2map\ snd\ snd\ snd\ (g\ x\ y) = dtor2\ y \rrbracket \Longrightarrow$
 $(\forall a \in JF1set\ z1. \forall x y. R2\ x\ y \wedge z1 = unfold1\ (case_prod\ f)\ (case_prod\ g)\ (x, y) \longrightarrow R1\ (fst\ a)\ (snd\ a)) \wedge$
 $(\forall a \in JF2set\ z2. \forall x y. R3\ x\ y \wedge z2 = unfold2\ (case_prod\ f)\ (case_prod\ g)\ (x, y) \longrightarrow R1\ (fst\ a)\ (snd\ a))$
 $\langle proof \rangle$

lemma $J_rel_coind_coind1$:

$\llbracket \forall x y. R2\ x\ y \longrightarrow (f\ x\ y \in F1in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))) \wedge$
 $F1map\ fst\ fst\ fst\ (f\ x\ y) = dtor1\ x \wedge$
 $F1map\ snd\ snd\ snd\ (f\ x\ y) = dtor1\ y;$
 $\forall x y. R3\ x\ y \longrightarrow (g\ x\ y \in F2in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))) \wedge$
 $F2map\ fst\ fst\ fst\ (g\ x\ y) = dtor2\ x \wedge$
 $F2map\ snd\ snd\ snd\ (g\ x\ y) = dtor2\ y \rrbracket \Longrightarrow$
 $(\exists y. R2\ x1\ y \wedge x1' = JF1map\ fst\ (unfold1\ (case_prod\ f)\ (case_prod\ g)\ (x1, y))) \longrightarrow x1' = x1) \wedge$
 $(\exists y. R3\ x2\ y \wedge x2' = JF2map\ fst\ (unfold2\ (case_prod\ f)\ (case_prod\ g)\ (x2, y))) \longrightarrow x2' = x2)$
 $\langle proof \rangle$

lemma $J_rel_coind_coind2$:

$\llbracket \forall x y. R2\ x\ y \longrightarrow (f\ x\ y \in F1in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))) \wedge$
 $F1map\ fst\ fst\ fst\ (f\ x\ y) = dtor1\ x \wedge$
 $F1map\ snd\ snd\ snd\ (f\ x\ y) = dtor1\ y;$
 $\forall x y. R3\ x\ y \longrightarrow (g\ x\ y \in F2in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))) \wedge$
 $F2map\ fst\ fst\ fst\ (g\ x\ y) = dtor2\ x \wedge$
 $F2map\ snd\ snd\ snd\ (g\ x\ y) = dtor2\ y \rrbracket \Longrightarrow$
 $(\exists x. R2\ x\ y1 \wedge y1' = JF1map\ snd\ (unfold1\ (case_prod\ f)\ (case_prod\ g)\ (x, y1))) \longrightarrow y1' = y1) \wedge$
 $(\exists x. R3\ x\ y2 \wedge y2' = JF2map\ snd\ (unfold2\ (case_prod\ f)\ (case_prod\ g)\ (x, y2))) \longrightarrow y2' = y2)$
 $\langle proof \rangle$

lemma J_rel_coind :

assumes $CIH1: \forall x2\ y2. R2\ x2\ y2 \longrightarrow F1rel\ R1\ R2\ R3\ (dtor1\ x2)\ (dtor1\ y2)$
and $CIH2: \forall x3\ y3. R3\ x3\ y3 \longrightarrow F2rel\ R1\ R2\ R3\ (dtor2\ x3)\ (dtor2\ y3)$
shows $R2 \leq JF1rel\ R1 \wedge R3 \leq JF2rel\ R1$
 $\langle proof \rangle$

lemma $JF1rel_F1rel$: $JF1rel\ R\ a\ b \longleftrightarrow F1rel\ R\ (JF1rel\ R)\ (JF2rel\ R)\ (dtor1\ a)\ (dtor1\ b)$

$\langle proof \rangle$

lemma $JF2rel_F2rel$: $JF2rel\ R\ a\ b \longleftrightarrow F2rel\ R\ (JF1rel\ R)\ (JF2rel\ R)\ (dtor2\ a)\ (dtor2\ b)$

$\langle proof \rangle$

lemma $JFrel_Comp_le$:

$JF1rel\ R1\ OO\ JF1rel\ R2 \leq JF1rel\ (R1\ OO\ R2) \wedge JF2rel\ R1\ OO\ JF2rel\ R2 \leq JF2rel\ (R1\ OO\ R2)$

$\langle proof \rangle$

context includes $lifting_syntax$

begin

lemma $unfold_transfer$:

$((S \Longrightarrow F1rel\ R\ S\ T) \Longrightarrow (T \Longrightarrow F2rel\ R\ S\ T) \Longrightarrow S \Longrightarrow JF1rel\ R)\ unfold1\ unfold1 \wedge$

$((S \Longrightarrow F1rel\ R\ S\ T) \Longrightarrow (T \Longrightarrow F2rel\ R\ S\ T) \Longrightarrow T \Longrightarrow JF2rel\ R)\ unfold2\ unfold2$

$\langle proof \rangle$

end

$\langle ML \rangle$

bnf $'a\ JF1$

$map: JF1map$

$sets: JF1set$

bd: *bd_F*
wits: *JF1wit*
rel: *JF1rel*
 ⟨*proof*⟩

bnf *'a JF2*
map: *JF2map*
sets: *JF2set*
bd: *bd_F*
wits: *JF2wit*
rel: *JF2rel*
 ⟨*proof*⟩

3 Normalized Composition of BNFs

Expected normal form: outer m-ary BNF is composed with m inner n-ary BNFs.

declare *[[bnf_internals]]*
bnf-axiomatization (*dead 'p1, F1set1: 'a1, F1set2: 'a2*) *F1*
 [*wits: ('p1, 'a1, 'a2) F1*]
for *map: F1map rel: F1rel*
bnf-axiomatization (*dead 'p2, F2set1: 'a1, F2set2: 'a2*) *F2*
 [*wits: 'a1 ⇒ ('p2, 'a1, 'a2) F2 'a2 ⇒ ('p2, 'a1, 'a2) F2*]
for *map: F2map rel: F2rel*
bnf-axiomatization (*dead 'p3, F3set1: 'a1, F3set2: 'a2*) *F3*
 [*wits: 'a1 ⇒ 'a2 ⇒ ('p3, 'a1, 'a2) F3*]
for *map: F3map rel: F3rel*
bnf-axiomatization (*dead 'p, Gset1: 'b1, Gset2: 'b2, Gset3: 'b3*) *G*
 [*wits: 'b1 ⇒ 'b3 ⇒ ('p, 'b1, 'b2, 'b3) G 'b2 ⇒ 'b3 ⇒ ('p, 'b1, 'b2, 'b3) G*]
for *map: Gmap rel: Grel*
type-synonym (*'p1, 'p2, 'p3, 'p, 'a1, 'a2*) *H =*
 (*'p, ('p1, 'a1, 'a2) F1, ('p2, 'a1, 'a2) F2, ('p3, 'a1, 'a2) F3*) *G*
type-synonym (*'p1, 'p2, 'p3, 'p*) *Hbd_type =*
 (*'p1 bd_type_F1 + 'p2 bd_type_F2 + 'p3 bd_type_F3*) × *'p bd_type_G*

abbreviation *F1in* **where** *F1in A1 A2 ≡ {x. F1set1 x ⊆ A1 ∧ F1set2 x ⊆ A2}*
abbreviation *F2in* **where** *F2in A1 A2 ≡ {x. F2set1 x ⊆ A1 ∧ F2set2 x ⊆ A2}*
abbreviation *F3in* **where** *F3in A1 A2 ≡ {x. F3set1 x ⊆ A1 ∧ F3set2 x ⊆ A2}*
abbreviation *Gin* **where** *Gin A1 A2 A3 ≡ {x. Gset1 x ⊆ A1 ∧ Gset2 x ⊆ A2 ∧ Gset3 x ⊆ A3}*

abbreviation *Gset* **where**
Gset ≡ BNF_Def.collect {Gset1, Gset2, Gset3}

abbreviation *Hmap* **::** (*'a1 ⇒ 'b1*) ⇒ (*'a2 ⇒ 'b2*) ⇒
 (*'p1, 'p2, 'p3, 'p, 'a1, 'a2*) *H ⇒ ('p1, 'p2, 'p3, 'p, 'b1, 'b2) H* **where**
Hmap f g ≡ Gmap (F1map f g) (F2map f g) (F3map f g)

abbreviation *Hset1* **::** (*'p1, 'p2, 'p3, 'p, 'a1, 'a2*) *H ⇒ 'a1 set* **where**
Hset1 ≡ Union o Gset o Gmap F1set1 F2set1 F3set1

abbreviation *Hset2* **::** (*'p1, 'p2, 'p3, 'p, 'a1, 'a2*) *H ⇒ 'a2 set* **where**
Hset2 ≡ Union o Gset o Gmap F1set2 F2set2 F3set2

lemma *Hset1_alt*:
Hset1 = Union o BNF_Def.collect {image F1set1 o Gset1, image F2set1 o Gset2, image F3set1 o Gset3}
 ⟨*proof*⟩

lemma *Hset2_alt*:
Hset2 = Union o BNF_Def.collect {image F1set2 o Gset1, image F2set2 o Gset2, image F3set2 o Gset3}
 ⟨*proof*⟩

abbreviation *Hbd* **where**

$Hbd \equiv (bd_F1 +c bd_F2 +c bd_F3) *c bd_G$

theorem *Hmap_id*: $Hmap\ id\ id = id$
 ⟨proof⟩

theorem *Hmap_comp*: $Hmap\ (f1\ o\ g1)\ (f2\ o\ g2) = Hmap\ f1\ f2\ o\ Hmap\ g1\ g2$
 ⟨proof⟩

theorem *Hmap_cong*: $\llbracket \bigwedge z. z \in Hset1\ x \implies f1\ z = g1\ z; \bigwedge z. z \in Hset2\ x \implies f2\ z = g2\ z \rrbracket \implies Hmap\ f1\ f2\ x = Hmap\ g1\ g2\ x$
 ⟨proof⟩

theorem *Hset1_natural*: $Hset1\ o\ Hmap\ f1\ f2 = image\ f1\ o\ Hset1$
 ⟨proof⟩

theorem *Hset2_natural*: $Hset2\ o\ Hmap\ f1\ f2 = image\ f2\ o\ Hset2$
 ⟨proof⟩

theorem *Hbd_card_order*: $card_order\ Hbd$
 ⟨proof⟩

theorem *Hbd_cinfinite*: $cinfinite\ Hbd$
 ⟨proof⟩

theorem *Hset1_bd*: $|Hset1\ (x :: ('p1, 'p2, 'p3, 'p, 'a1, 'a2)\ H)| \leq o$
 $(Hbd :: ('p1, 'p2, 'p3, 'p)\ Hbd_type\ rel)$
 ⟨proof⟩

theorem *Hset2_bd*: $|Hset2\ (x :: ('p1, 'p2, 'p3, 'p, 'a1, 'a2)\ H)| \leq o$
 $(Hbd :: ('p1, 'p2, 'p3, 'p)\ Hbd_type\ rel)$
 ⟨proof⟩

abbreviation *Hin where* $Hin\ A1\ A2 \equiv \{x. Hset1\ x \subseteq A1 \wedge Hset2\ x \subseteq A2\}$

lemma *Hin_alt*: $Hin\ A1\ A2 = Gin\ (F1in\ A1\ A2)\ (F2in\ A1\ A2)\ (F3in\ A1\ A2)$
 ⟨proof⟩

definition *Hwit1 where* $Hwit1\ b\ c = wit1_G\ wit_F1\ (wit_F3\ b\ c)$

definition *Hwit21 where* $Hwit21\ b\ c = wit2_G\ (wit1_F2\ b)\ (wit_F3\ b\ c)$

definition *Hwit22 where* $Hwit22\ b\ c = wit2_G\ (wit2_F2\ c)\ (wit_F3\ b\ c)$

lemma *Hwit1*:

$\bigwedge x. x \in Hset1\ (Hwit1\ b\ c) \implies x = b$
 $\bigwedge x. x \in Hset2\ (Hwit1\ b\ c) \implies x = c$
 ⟨proof⟩

lemma *Hwit21*:

$\bigwedge x. x \in Hset1\ (Hwit21\ b\ c) \implies x = b$
 $\bigwedge x. x \in Hset2\ (Hwit21\ b\ c) \implies x = c$
 ⟨proof⟩

lemma *Hwit22*:

$\bigwedge x. x \in Hset1\ (Hwit22\ b\ c) \implies x = b$
 $\bigwedge x. x \in Hset2\ (Hwit22\ b\ c) \implies x = c$
 ⟨proof⟩

lemma *Grel_cong*: $\llbracket R1 = S1; R2 = S2; R3 = S3 \rrbracket \implies Grel\ R1\ R2\ R3 = Grel\ S1\ S2\ S3$
 ⟨proof⟩

definition *Hrel where*

$Hrel\ R1\ R2 = (BNF_Def.Grp\ (Hin\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))))\ (Hmap\ fst\ fst)^{--1}\ OO$

$(BNF_Def.Grp\ (Hin\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))))\ (Hmap\ snd\ snd)$

lemmas $Hrel_unfold = trans[OF\ Hrel_def\ trans[OF\ OO_Grp_cong[OF\ Hin_alt]\ trans[OF\ arg_cong2[of\ __ __ relcompp, OF\ trans[OF\ arg_cong[of\ __ __ conversep, OF\ sym[OF\ G.rel_Grp]]\ G.rel_conversep[symmetric]]\ sym[OF\ G.rel_Grp]]\ trans[OF\ G.rel_compp[symmetric]\ Grel_cong[OF\ sym[OF\ F1.rel_compp_Grp]\ sym[OF\ F2.rel_compp_Grp]\ sym[OF\ F3.rel_compp_Grp]]]]]$

bnf $H: ('p1, 'p2, 'p3, 'p, 'a1, 'a2)\ H$
 $map: Hmap$
 $sets: Hset1\ Hset2$
 $bd: Hbd :: ('p1, 'p2, 'p3, 'p)\ Hbd_type\ rel$
 $rel: Hrel$
 $\langle proof \rangle$

4 Removing Live Variables

declare $[[bnf_internals]]$

bnf-axiomatization $(dead\ 'p, Fset1: 'a1, Fset2: 'a2, Fset3: 'a3)\ F\ \mathbf{for}\ map: Fmap\ rel: Frel$

abbreviation $F1map :: ('a2 \Rightarrow 'b2) \Rightarrow ('a3 \Rightarrow 'b3) \Rightarrow ('p, 'a1, 'a2, 'a3)\ F \Rightarrow ('p, 'a1, 'b2, 'b3)\ F\ \mathbf{where}$
 $F1map \equiv Fmap\ id$

abbreviation $F2map :: ('a3 \Rightarrow 'b3) \Rightarrow ('p, 'a1, 'a2, 'a3)\ F \Rightarrow ('p, 'a1, 'a2, 'b3)\ F\ \mathbf{where}$
 $F2map \equiv Fmap\ id\ id$

abbreviation $F1set1 \equiv Fset2$

abbreviation $F1set2 \equiv Fset3$

abbreviation $F2set \equiv Fset3$

theorem $F1map_id: F1map\ id\ id = id$
 $\langle proof \rangle$

theorem $F2map_id: F2map\ id = id$
 $\langle proof \rangle$

theorem $F1map_comp: F1map\ (f1\ o\ g1)\ (f2\ o\ g2) = F1map\ f1\ f2\ o\ F1map\ g1\ g2$
 $\langle proof \rangle$

theorem $F2map_comp: F2map\ (f\ o\ g) = F2map\ f\ o\ F2map\ g$
 $\langle proof \rangle$

theorem $F1map_cong: [\bigwedge z. z \in F1set1\ x \implies f1\ z = g1\ z; \bigwedge z. z \in F1set2\ x \implies f2\ z = g2\ z] \implies F1map\ f1\ f2\ x = F1map\ g1\ g2\ x$
 $\langle proof \rangle$

theorem $F2map_cong: [\bigwedge z. z \in F2set\ x \implies f\ z = g\ z] \implies F2map\ f\ x = F2map\ g\ x$
 $\langle proof \rangle$

theorem $F1set1_natural: F1set1\ o\ F1map\ f1\ f2 = image\ f1\ o\ F1set1$
 $\langle proof \rangle$

theorem $F1set2_natural: F1set2\ o\ F1map\ f1\ f2 = image\ f2\ o\ F1set2$
 $\langle proof \rangle$

theorem $F2set_natural: F2set\ o\ F2map\ f = image\ f\ o\ F2set$
 $\langle proof \rangle$

abbreviation $Fin :: 'a1\ set \Rightarrow 'a2\ set \Rightarrow 'a3\ set \Rightarrow ((('p, 'a1, 'a2, 'a3)\ F)\ set)\ \mathbf{where}$
 $Fin\ A1\ A2\ A3 \equiv \{x. Fset1\ x \subseteq A1 \wedge Fset2\ x \subseteq A2 \wedge Fset3\ x \subseteq A3\}$

abbreviation $F1in :: 'a2 \text{ set} \Rightarrow 'a3 \text{ set} \Rightarrow (('p, 'a1, 'a2, 'a3) F) \text{ set}$ **where**
 $F1in A1 A2 \equiv \{x. F1set1 x \subseteq A1 \wedge F1set2 x \subseteq A2\}$

lemma $F1in_alt: F1in A2 A3 = Fin UNIV A2 A3$
 $\langle proof \rangle$

abbreviation $F2in :: 'a3 \text{ set} \Rightarrow (('p, 'a1, 'a2, 'a3) F) \text{ set}$ **where**
 $F2in A \equiv \{x. F2set x \subseteq A\}$

lemma $F2in_alt: F2in A3 = Fin UNIV UNIV A3$
 $\langle proof \rangle$

lemma $Frel_cong: [R1 = S1; R2 = S2; R3 = S3] \Longrightarrow Frel R1 R2 R3 = Frel S1 S2 S3$
 $\langle proof \rangle$

definition $F1rel$ **where**

$F1rel R1 R2 = (BNF_Def.Grp (F1in (Collect (case_prod R1))) (Collect (case_prod R2))) (F1map fst fst) \hat{\ } --1$
 OO
 $(BNF_Def.Grp (F1in (Collect (case_prod R1))) (Collect (case_prod R2))) (F1map snd snd)$

lemmas $F1rel_unfold = trans[OF F1rel_def trans[OF OO_Grp_cong[OF F1in_alt]$
 $trans[OF arg_cong2[of _ _ _ relcompp, OF trans[OF arg_cong[of _ _ conversep, OF sym[OF F.rel_Grp]]$
 $F.rel_conversep[symmetric]] sym[OF F.rel_Grp]]$
 $trans[OF F.rel_compp[symmetric] Frel_cong[OF trans[OF Grp_UNIV_id[OF refl] eq_alt[symmetric]]$
 $Grp_fst_snd Grp_fst_snd]]]]]$

definition $F2rel$ **where**

$F2rel R1 = (BNF_Def.Grp (F2in (Collect (case_prod R1))) (F2map fst)) \hat{\ } --1 OO$
 $(BNF_Def.Grp (F2in (Collect (case_prod R1))) (F2map snd))$

lemmas $F2rel_unfold = trans[OF F2rel_def trans[OF OO_Grp_cong[OF F2in_alt]$
 $trans[OF arg_cong2[of _ _ _ relcompp, OF trans[OF arg_cong[of _ _ conversep, OF sym[OF F.rel_Grp]]$
 $F.rel_conversep[symmetric]] sym[OF F.rel_Grp]]$
 $trans[OF F.rel_compp[symmetric] Frel_cong[OF trans[OF Grp_UNIV_id[OF refl] eq_alt[symmetric]]$
 $trans[OF Grp_UNIV_id[OF refl] eq_alt[symmetric]] Grp_fst_snd]]]]]$

bnf $F1: ('p, 'a1, 'a2, 'a3) F$
 $map: F1map$
 $sets: F1set1 F1set2$
 $bd: bd_F :: ('p bd_type_F) rel$
 $rel: F1rel$
 $\langle proof \rangle$

bnf $F2: ('p, 'a1, 'a2, 'a3) F$
 $map: F2map$
 $sets: F2set$
 $bd: bd_F :: ('p bd_type_F) rel$
 $rel: F2rel$
 $\langle proof \rangle$

5 Adding New Live Variables

declare $[[bnf_internals]]$

bnf-axiomatization $(dead 'p, Fset1: 'a1, Fset2: 'a2) F$
 $[wits: 'a1 \Rightarrow 'a2 \Rightarrow ('p, 'a1, 'a2) F]$
for $map: Fmap rel: Frel$

type-synonym $('p, 'a1, 'a2, 'a3, 'a4) F' = ('p, 'a3, 'a4) F$

abbreviation $F'map :: ('a1 \Rightarrow 'b1) \Rightarrow ('a2 \Rightarrow 'b2) \Rightarrow ('a3 \Rightarrow 'b3) \Rightarrow ('a4 \Rightarrow 'b4) \Rightarrow ('p, 'a1, 'a2, 'a3, 'a4) F'$
 $\Rightarrow ('p, 'b1, 'b2, 'b3, 'b4) F'$ **where**
 $F'map f1 f2 f3 f4 \equiv Fmap f3 f4$

abbreviation $F'set1 :: ('p, 'a1, 'a2, 'a3, 'a4) F' \Rightarrow 'a1 \text{ set}$ **where**
 $F'set1 \equiv \lambda_ . \{ \}$

abbreviation $F'set2 :: ('p, 'a1, 'a2, 'a3, 'a4) F' \Rightarrow 'a2 \text{ set}$ **where**
 $F'set2 \equiv \lambda_ . \{ \}$

abbreviation $F'set3 :: ('p, 'a1, 'a2, 'a3, 'a4) F' \Rightarrow 'a3 \text{ set}$ **where**
 $F'set3 \equiv F'set1$

abbreviation $F'set4 :: ('p, 'a1, 'a2, 'a3, 'a4) F' \Rightarrow 'a4 \text{ set}$ **where**
 $F'set4 \equiv F'set2$

abbreviation $F'bd$ **where**
 $F'bd \equiv bd_F$

theorem $F'map_id$: $F'map \text{ id id id id id} = \text{id}$
 $\langle \text{proof} \rangle$

theorem $F'map_comp$:
 $F'map (f1 \circ g1) (f2 \circ g2) (f3 \circ g3) (f4 \circ g4) = F'map f1 f2 f3 f4 \circ F'map g1 g2 g3 g4$
 $\langle \text{proof} \rangle$

theorem $F'map_cong$:
 $\llbracket \bigwedge z. z \in F'set1 x \implies f1 z = g1 z; \bigwedge z. z \in F'set2 x \implies f2 z = g2 z;$
 $\bigwedge z. z \in F'set3 x \implies f3 z = g3 z; \bigwedge z. z \in F'set4 x \implies f4 z = g4 z \rrbracket$
 $\implies F'map f1 f2 f3 f4 x = F'map g1 g2 g3 g4 x$
 $\langle \text{proof} \rangle$

theorem $F'set1_natural$: $F'set1 \circ F'map f1 f2 f3 f4 = \text{image } f1 \circ F'set1$
 $\langle \text{proof} \rangle$

theorem $F'set2_natural$: $F'set2 \circ F'map f1 f2 f3 f4 = \text{image } f2 \circ F'set2$
 $\langle \text{proof} \rangle$

theorem $F'set3_natural$: $F'set3 \circ F'map f1 f2 f3 f4 = \text{image } f3 \circ F'set3$
 $\langle \text{proof} \rangle$

theorem $F'set4_natural$: $F'set4 \circ F'map f1 f2 f3 f4 = \text{image } f4 \circ F'set4$
 $\langle \text{proof} \rangle$

theorem $F'bd_card_order$: $\text{card_order } bd_F$
 $\langle \text{proof} \rangle$

theorem $F'bd_cinfinte$: $\text{cinfinte } bd_F$
 $\langle \text{proof} \rangle$

theorem $F'set1_bd$: $|F'set1 x| \leq o F'bd$
 $\langle \text{proof} \rangle$

theorem $F'set2_bd$: $|F'set2 x| \leq o F'bd$
 $\langle \text{proof} \rangle$

theorem $F'set3_bd$: $|F'set3 (x :: ('c, 'a, 'd) F)| \leq o (F'bd :: 'c \text{ bd_type_} F \text{ rel})$
 $\langle \text{proof} \rangle$

theorem $F'set4_bd$: $|F'set4 (x :: ('c, 'a, 'd) F)| \leq o (F'bd :: 'c \text{ bd_type_} F \text{ rel})$
 $\langle \text{proof} \rangle$

abbreviation $F'in :: 'a1 \text{ set} \Rightarrow 'a2 \text{ set} \Rightarrow 'a3 \text{ set} \Rightarrow 'a4 \text{ set} \Rightarrow (('p, 'a1, 'a2, 'a3, 'a4) F') \text{ set}$ **where**
 $F'in A1 A2 A3 A4 \equiv \{ x. F'set1 x \subseteq A1 \wedge F'set2 x \subseteq A2 \wedge F'set3 x \subseteq A3 \wedge F'set4 x \subseteq A4 \}$

definition $F'rel$ **where**

$F' \text{rel } R1 \ R2 \ R3 \ R4 = (\text{BNF_Def.Grp } (F' \text{in } (\text{Collect } (\text{case_prod } R1)) (\text{Collect } (\text{case_prod } R2)) (\text{Collect } (\text{case_prod } R3)) (\text{Collect } (\text{case_prod } R4)))) (F' \text{map } \text{fst } \text{fst } \text{fst } \text{fst}) \wedge \text{---} 1 \text{ OO}$
 $(\text{BNF_Def.Grp } (F' \text{in } (\text{Collect } (\text{case_prod } R1)) (\text{Collect } (\text{case_prod } R2)) (\text{Collect } (\text{case_prod } R3)) (\text{Collect } (\text{case_prod } R4)))) (F' \text{map } \text{snd } \text{snd } \text{snd } \text{snd})$

lemmas $F' \text{rel_unfold} = \text{trans}[\text{OF } F' \text{rel_def}[\text{unfolded } \text{eqTrueI}[\text{OF } \text{empty_subsetI}] \text{ simp_thms}(22)]]$
 $\text{trans}[\text{OF } \text{OO_Grp_cong}[\text{OF } \text{refl}] \text{ sym}[\text{OF } F' \text{rel_compp_Grp}]]]$

bnf F' : $(p, a1, a2, a3, a4) F'$
 $\text{map: } F' \text{map}$
 $\text{sets: } F' \text{set1 } F' \text{set2 } F' \text{set3 } F' \text{set4}$
 $\text{bd: } F' \text{bd} :: p \text{ bd_type_} F \text{ rel}$
 $\text{wits: } \text{wit_} F$
 $\text{rel: } F' \text{rel}$
 $\text{plugins del: lifting transfer}$
 $\langle \text{proof} \rangle$

6 Changing the Order of Live Variables

declare $[[\text{bnf_internals}]]$

bnf-axiomatization $(\text{dead } p, F \text{set1: } a1, F \text{set2: } a2, F \text{set3: } a3) F$ **for** $\text{map: } F \text{map } \text{rel: } F \text{rel}$

type-synonym $(p, a1, a2, a3) F' = (p, a3, a1, a2) F$

abbreviation $Fin :: a1 \text{ set} \Rightarrow a2 \text{ set} \Rightarrow a3 \text{ set} \Rightarrow ((p, a1, a2, a3) F) \text{ set}$ **where**
 $Fin \ A1 \ A2 \ A3 \equiv \{x. F \text{set1 } x \subseteq A1 \wedge F \text{set2 } x \subseteq A2 \wedge F \text{set3 } x \subseteq A3\}$

abbreviation $F' \text{map} :: (a1 \Rightarrow b1) \Rightarrow (a2 \Rightarrow b2) \Rightarrow (a3 \Rightarrow b3) \Rightarrow (p, a1, a2, a3) F' \Rightarrow (p, b1, b2, b3) F'$ **where**
 $F' \text{map } f \ g \ h \equiv F \text{map } h \ f \ g$

abbreviation $F' \text{set1} :: (p, a1, a2, a3) F' \Rightarrow a1 \text{ set}$ **where**
 $F' \text{set1} \equiv F \text{set2}$

abbreviation $F' \text{set2} :: (p, a1, a2, a3) F' \Rightarrow a2 \text{ set}$ **where**
 $F' \text{set2} \equiv F \text{set3}$

abbreviation $F' \text{set3} :: (p, a1, a2, a3) F' \Rightarrow a3 \text{ set}$ **where**
 $F' \text{set3} \equiv F \text{set1}$

abbreviation $F' \text{bd}$ **where**
 $F' \text{bd} \equiv \text{bd_} F$

theorem $F' \text{map_id: } F' \text{map } \text{id } \text{id } \text{id} = \text{id}$
 $\langle \text{proof} \rangle$

theorem $F' \text{map_comp: } F' \text{map } (f1 \circ g1) (f2 \circ g2) (f3 \circ g3) = F' \text{map } f1 \ f2 \ f3 \circ F' \text{map } g1 \ g2 \ g3$
 $\langle \text{proof} \rangle$

theorem $F' \text{map_cong: } [[\bigwedge z. z \in F' \text{set1 } x \Longrightarrow f1 \ z = g1 \ z; \bigwedge z. z \in F' \text{set2 } x \Longrightarrow f2 \ z = g2 \ z; \bigwedge z. z \in F' \text{set3 } x \Longrightarrow f3 \ z = g3 \ z]]$
 $\Longrightarrow F' \text{map } f1 \ f2 \ f3 \ x = F' \text{map } g1 \ g2 \ g3 \ x$
 $\langle \text{proof} \rangle$

theorem $F' \text{set1_natural: } F' \text{set1 } \circ F' \text{map } f1 \ f2 \ f3 = \text{image } f1 \ \circ F' \text{set1}$
 $\langle \text{proof} \rangle$

theorem $F' \text{set2_natural: } F' \text{set2 } \circ F' \text{map } f1 \ f2 \ f3 = \text{image } f2 \ \circ F' \text{set2}$
 $\langle \text{proof} \rangle$

theorem $F' \text{set3_natural: } F' \text{set3 } \circ F' \text{map } f1 \ f2 \ f3 = \text{image } f3 \ \circ F' \text{set3}$
 $\langle \text{proof} \rangle$

theorem $F'bd_card_order$: $card_order\ F'bd$
 ⟨proof⟩

theorem $F'bd_cinfinte$: $cinfinte\ F'bd$
 ⟨proof⟩

theorem $F'set1_bd$: $|F'set1\ (x :: ('c, 'a, 'b, 'd)\ F)| \leq o\ (F'bd :: 'c\ bd_type_F\ rel)$
 ⟨proof⟩

theorem $F'set2_bd$: $|F'set2\ (x :: ('c, 'a, 'b, 'd)\ F)| \leq o\ (F'bd :: 'c\ bd_type_F\ rel)$
 ⟨proof⟩

theorem $F'set3_bd$: $|F'set3\ (x :: ('c, 'a, 'b, 'd)\ F)| \leq o\ (F'bd :: 'c\ bd_type_F\ rel)$
 ⟨proof⟩

abbreviation $F'in :: 'a1\ set \Rightarrow 'a2\ set \Rightarrow 'a3\ set \Rightarrow (('p, 'a1, 'a2, 'a3)\ F')\ set$ **where**
 $F'in\ A1\ A2\ A3 \equiv \{x. F'set1\ x \subseteq A1 \wedge F'set2\ x \subseteq A2 \wedge F'set3\ x \subseteq A3\}$

lemma $F'in_alt$: $F'in\ A1\ A2\ A3 = Fin\ A3\ A1\ A2$
 ⟨proof⟩

definition $F'rel$ **where**
 $F'rel\ R1\ R2\ R3 = (BNF_Def.Grp\ (F'in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))))\ (F'map\ fst\ fst\ fst) \hat{\ }_{--1}\ OO$
 $(BNF_Def.Grp\ (F'in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))))\ (F'map\ snd\ snd\ snd)$

lemmas $F'rel_unfold = trans[OF\ F'rel_def\ trans[OF\ OO_Grp_cong[OF\ F'in_alt]\ sym[OF\ F.rel_compp_Grp]]]$

bnf F' : $('p, 'a1, 'a2, 'a3)\ F'$
 map : $F'map$
 $sets$: $F'set1\ F'set2\ F'set3$
 bd : $F'bd :: 'p\ bd_type_F\ rel$
 rel : $F'rel$
 ⟨proof⟩

7 Mutual View on Nested Datatypes

notation $BNF_Def.convolve\ (<_, _>)$

declare $[[bnf_internals]]$

declare $[[typedef_overloaded]]$

bnf-axiomatization $('a, 'b)\ F0$ $[wits: 'a \Rightarrow ('a, 'b)\ F0]$
bnf-axiomatization $('a, 'b)\ G0$ $[wits: 'a \Rightarrow ('a, 'b)\ G0]$

7.1 Nested Definition

datatype $'a\ F = CF\ ('a, 'a\ F)\ F0$
datatype $'a\ G = CG\ ('a, ('a\ G)\ F)\ G0$

type-synonym $('b, 'c)\ F_pre_F = ('c, 'b)\ F0$
type-synonym $('c, 'a)\ G_pre_G = ('a, 'c)\ F0$

term $ctor_fold_F :: (('b, 'c)\ F_pre_F \Rightarrow 'b) \Rightarrow 'c\ F \Rightarrow 'b$
term $ctor_fold_G :: (('c, 'a)\ G_pre_G \Rightarrow 'c) \Rightarrow 'a\ G \Rightarrow 'c$
term $ctor_rec_F :: (('c\ F \times 'b, 'c)\ F_pre_F \Rightarrow 'b) \Rightarrow 'c\ F \Rightarrow 'b$
term $ctor_rec_G :: (('a\ G \times 'c, 'a)\ G_pre_G \Rightarrow 'c) \Rightarrow 'a\ G \Rightarrow 'c$
thm $F.ctor_rel_induct$
thm $G.ctor_rel_induct[unfolded\ rel_pre_G_def\ id_apply]$

7.2 Isomorphic Mutual Definition

datatype $'a$ $G_M = CG ('a, 'a GF_M) G0$
and $'a$ $GF_M = CF ('a G_M, 'a GF_M) F0$

type-synonym $('b, 'c) GF_{M_pre_}GF_M = ('c, 'b) F0$
type-synonym $('c, 'a) G_{M_pre_}G_M = ('a, 'c) G0$

term $ctor_fold_G_M :: (('c, 'a) G_{M_pre_}G_M \Rightarrow 'b) \Rightarrow (('c, 'b) GF_{M_pre_}GF_M \Rightarrow 'c) \Rightarrow 'a G_M \Rightarrow 'b$
term $ctor_fold_GF_M :: (('c, 'a) G_{M_pre_}G_M \Rightarrow 'b) \Rightarrow (('c, 'b) GF_{M_pre_}GF_M \Rightarrow 'c) \Rightarrow 'a GF_M \Rightarrow 'c$
term $ctor_rec_G_M :: (('a GF_M \times 'c, 'a) G_{M_pre_}G_M \Rightarrow 'b) \Rightarrow (('a GF_M \times 'c, 'a G_M \times 'b) GF_{M_pre_}GF_M \Rightarrow 'c) \Rightarrow 'a G_M \Rightarrow 'b$
term $ctor_rec_GF_M :: (('a GF_M \times 'c, 'a) G_{M_pre_}G_M \Rightarrow 'b) \Rightarrow (('a GF_M \times 'c, 'a G_M \times 'b) GF_{M_pre_}GF_M \Rightarrow 'c) \Rightarrow 'a GF_M \Rightarrow 'c$
thm $G_M_GF_M.ctor_rel_induct[unfolded rel_pre_G_M_def rel_pre_GF_M_def]$

7.3 Mutualization

7.3.1 Iterators

definition $n2m_ctor_fold_G :: (('c, 'a) G_{M_pre_}G_M \Rightarrow 'b) \Rightarrow (('c, 'b) GF_{M_pre_}GF_M \Rightarrow 'c) \Rightarrow 'a G \Rightarrow 'b$
where $n2m_ctor_fold_G s1 s2 = ctor_fold_G (s1 o$
 $map_pre_G_M id (id :: unit \Rightarrow unit) (ctor_fold_F (s2 o BNF_Composition.id_bnf o BNF_Composition.id_bnf)))$
 $o BNF_Composition.id_bnf o BNF_Composition.id_bnf)$
definition $n2m_ctor_fold_G_F :: (('c, 'a) G_{M_pre_}G_M \Rightarrow 'b) \Rightarrow (('c, 'b) GF_{M_pre_}GF_M \Rightarrow 'c) \Rightarrow 'a G F \Rightarrow 'c$
where $n2m_ctor_fold_G_F s1 s2 = ctor_fold_F (s2 o map_pre_GF_M (id :: unit \Rightarrow unit) (n2m_ctor_fold_G$
 $s1 s2) id o BNF_Composition.id_bnf o BNF_Composition.id_bnf)$

lemma $G_ctor_o_fold: ctor_fold_G s o ctor_G = s o map_pre_G id (ctor_fold_G s)$
 $\langle proof \rangle$

lemma $F_ctor_o_fold: ctor_fold_F s o ctor_F = s o map_pre_F id (ctor_fold_F s)$
 $\langle proof \rangle$

lemma $G_ctor_o_rec: ctor_rec_G s o ctor_G = s o map_pre_G id (BNF_Def.convolve id (ctor_rec_G s))$
 $\langle proof \rangle$

lemma $F_ctor_o_rec: ctor_rec_F s o ctor_F = s o map_pre_F id (BNF_Def.convolve id (ctor_rec_F s))$
 $\langle proof \rangle$

lemma $n2m_ctor_fold_G:$
 $n2m_ctor_fold_G s1 s2 o ctor_G = s1 o map_pre_G_M id id (n2m_ctor_fold_G_F s1 s2) o BNF_Composition.id_bnf$
 $o BNF_Composition.id_bnf$
 $\langle proof \rangle$

lemma $n2m_ctor_fold_G_F:$
 $n2m_ctor_fold_G_F s1 s2 o ctor_F = s2 o map_pre_GF_M id (n2m_ctor_fold_G s1 s2) (n2m_ctor_fold_G_F$
 $s1 s2) o BNF_Composition.id_bnf o BNF_Composition.id_bnf$
 $\langle proof \rangle$

7.3.2 Recursors

definition $n2m_ctor_rec_G ::$
 $(('a G F \times 'c, 'a) G_{M_pre_}G_M \Rightarrow 'b) \Rightarrow (('a G F \times 'c, 'a G \times 'b) GF_{M_pre_}GF_M \Rightarrow 'c) \Rightarrow 'a G \Rightarrow 'b$
where $n2m_ctor_rec_G s1 s2 =$
 $ctor_rec_G (s1 o$
 $map_pre_G_M id (id :: unit \Rightarrow unit)$
 $(BNF_Def.convolve (map_F fst) (ctor_rec_F (s2 o map_pre_GF_M (id :: unit \Rightarrow unit) id (map_prod (map_F$
 $fst) id) o BNF_Composition.id_bnf o BNF_Composition.id_bnf))) o$
 $BNF_Composition.id_bnf o BNF_Composition.id_bnf)$

definition $n2m_ctor_rec_G_F ::$
 $(('a G F \times 'c, 'a) G_{M_pre_}G_M \Rightarrow 'b) \Rightarrow (('a G F \times 'c, 'a G \times 'b) GF_{M_pre_}GF_M \Rightarrow 'c) \Rightarrow 'a G F \Rightarrow 'c$
where $n2m_ctor_rec_G_F s1 s2 = ctor_rec_F (s2 o map_pre_GF_M (id :: unit \Rightarrow unit) (BNF_Def.convolve id$
 $(n2m_ctor_rec_G s1 s2)) id o BNF_Composition.id_bnf o BNF_Composition.id_bnf)$

lemma $n2m_ctor_rec_G$:
 $n2m_ctor_rec_G\ s1\ s2\ o\ ctor_G = s1\ o\ map_pre_G_M\ id\ id\ (BNF_Def.convolve\ id\ (n2m_ctor_rec_G_F\ s1\ s2))$
 $o\ BNF_Composition.id_bnf\ o\ BNF_Composition.id_bnf$
 $\langle proof \rangle$

lemma $n2m_ctor_rec_G_F$:
 $n2m_ctor_rec_G_F\ s1\ s2\ o\ ctor_F = s2\ o\ map_pre_GF_M\ id\ (BNF_Def.convolve\ id\ (n2m_ctor_rec_G\ s1\ s2))$
 $(BNF_Def.convolve\ id\ (n2m_ctor_rec_G_F\ s1\ s2))\ o\ BNF_Composition.id_bnf\ o\ BNF_Composition.id_bnf$
 $\langle proof \rangle$

7.3.3 Induction

lemma $n2m_rel_induct_G_G_F$:
assumes $IH1: \forall x\ y. BNF_Def.vimage2p\ (BNF_Composition.id_bnf\ o\ BNF_Composition.id_bnf)\ (BNF_Composition.id_bnf\ o\ BNF_Composition.id_bnf)\ (rel_pre_G_M\ P\ R\ S)\ x\ y \longrightarrow R\ (ctor_G\ x)\ (ctor_G\ y)$
and $IH2: \forall x\ y. BNF_Def.vimage2p\ (BNF_Composition.id_bnf\ o\ BNF_Composition.id_bnf)\ (BNF_Composition.id_bnf\ o\ BNF_Composition.id_bnf)\ (rel_pre_GF_M\ P\ R\ S)\ x\ y \longrightarrow S\ (ctor_F\ x)\ (ctor_F\ y)$
shows $rel_G\ P \leq R \wedge rel_F\ (rel_G\ P) \leq S$
 $\langle proof \rangle$

lemmas $n2m_ctor_induct_G_G_F = spec[OF\ spec\ [OF\ n2m_rel_induct_G_G_F\ of\ (=)\ BNF_Def.Grp\ (Collect\ R)\ id\ BNF_Def.Grp\ (Collect\ S)\ id\ for\ R\ S,$
 $unfolded\ G.rel_eq\ F.rel_eq\ eq_le_Grp_id_iff\ all_simps(1,2)[symmetric]]],$
 $unfolded\ eq_alt\ pre_G_M.rel_Grp\ pre_GF_M.rel_Grp\ pre_G_M.map_id0\ pre_GF_M.map_id0,$
 $unfolded\ vimage2p_comp\ vimage2p_id\ comp_apply\ comp_id\ Grp_id_mono_subst$
 $type_copy_vimage2p_Grp_Rep[OF\ BNF_Composition.type_definition_id_bnf_UNIV]$
 $type_copy_Abs_o_Rep[OF\ BNF_Composition.type_definition_id_bnf_UNIV]$
 $eqTrueI[OF\ subset_UNIV]\ simp_thms(22)$
 $atomize_conjL[symmetric]\ atomize_all[symmetric]\ atomize_imp[symmetric],$
 $unfolded\ subset_iff\ mem_Collect_eq]$

8 Mutual View on Nested Coatypes

bnf-axiomatization $(a, b)\ coF0$
bnf-axiomatization $(a, b)\ coG0$

8.1 Nested definition

codatatype $a\ coF = CcoF\ (a, a\ coF)\ coF0$
codatatype $a\ coG = CcoG\ (a, (a\ coG)\ coF)\ coG0$

type-synonym $(b, c)\ coF_pre_coF = (c, b)\ coF0$
type-synonym $(c, a)\ coG_pre_coG = (a, c\ coF)\ coG0$

term $dtor_unfold_coF :: (b \Rightarrow (b, c)\ coF_pre_coF) \Rightarrow b \Rightarrow c\ coF$
term $dtor_unfold_coG :: (c \Rightarrow (c, a)\ coG_pre_coG) \Rightarrow c \Rightarrow a\ coG$
term $dtor_corec_coF :: (b \Rightarrow (c\ coF + b, c)\ coF_pre_coF) \Rightarrow b \Rightarrow c\ coF$
term $dtor_corec_coG :: (c \Rightarrow (a\ coG + c, a)\ coG_pre_coG) \Rightarrow c \Rightarrow a\ coG$
thm $coF.dtor_rel_coinduct$
thm $coG.dtor_rel_coinduct[unfolded\ rel_pre_coG_def\ id_apply]$

8.2 Isomorphic Mutual Definition

codatatype $a\ coG_M = CcoG\ (a, a\ coGcoF_M)\ coG0$
and $a\ coGcoF_M = CcoF\ (a\ coG_M, a\ coGcoF_M)\ coF0$

type-synonym $(b, c)\ coGcoF_M_pre_coGcoF_M = (c, b)\ coF0$
type-synonym $(c, a)\ coG_M_pre_coG_M = (a, c)\ coG0$

term $dtor_unfold_coG_M :: (b \Rightarrow (c, a)\ coG_M_pre_coG_M) \Rightarrow (c \Rightarrow (c, b)\ coGcoF_M_pre_coGcoF_M) \Rightarrow b \Rightarrow a\ coG_M$
term $dtor_unfold_coGcoF_M :: (b \Rightarrow (c, a)\ coG_M_pre_coG_M) \Rightarrow (c \Rightarrow (c, b)\ coGcoF_M_pre_coGcoF_M) \Rightarrow c \Rightarrow a\ coGcoF_M$

term $\text{dtor_corec_coG}_M :: ('b \Rightarrow ('a \text{ coGcoF}_M + 'c, 'a) \text{ coG}_M\text{_pre_coG}_M) \Rightarrow ('c \Rightarrow ('a \text{ coGcoF}_M + 'c, 'a \text{ coG}_M + 'b) \text{ coGcoF}_M\text{_pre_coGcoF}_M) \Rightarrow 'b \Rightarrow 'a \text{ coG}_M$
term $\text{dtor_corec_coGcoF}_M :: ('b \Rightarrow ('a \text{ coGcoF}_M + 'c, 'a) \text{ coG}_M\text{_pre_coG}_M) \Rightarrow ('c \Rightarrow ('a \text{ coGcoF}_M + 'c, 'a \text{ coG}_M + 'b) \text{ coGcoF}_M\text{_pre_coGcoF}_M) \Rightarrow 'c \Rightarrow 'a \text{ coGcoF}_M$
thm $\text{coG}_M\text{_coGcoF}_M.\text{dtor_rel_coinduct}[\text{unfolded_rel_pre_coG}_M\text{_def_rel_pre_coGcoF}_M\text{_def}]$

8.3 Mutualization

8.3.1 Coiterators

definition $\text{n2m_dtor_unfold_coG} :: ('b \Rightarrow ('c, 'a) \text{ coG}_M\text{_pre_coG}_M) \Rightarrow ('c \Rightarrow ('c, 'b) \text{ coGcoF}_M\text{_pre_coGcoF}_M) \Rightarrow 'b \Rightarrow 'a \text{ coG}$

where $\text{n2m_dtor_unfold_coG } s1 \ s2 = \text{dtor_unfold_coG } (\text{BNF_Composition.id_bnf } o \ \text{BNF_Composition.id_bnf } o \ \text{map_pre_coG}_M \ \text{id } (\text{id} :: \text{unit} \Rightarrow \text{unit}) \ (\text{dtor_unfold_coF } (\text{BNF_Composition.id_bnf } o \ \text{BNF_Composition.id_bnf } o \ s2))) \ o \ s1$

definition $\text{n2m_dtor_unfold_coG_coF} :: ('b \Rightarrow ('c, 'a) \text{ coG}_M\text{_pre_coG}_M) \Rightarrow ('c \Rightarrow ('c, 'b) \text{ coGcoF}_M\text{_pre_coGcoF}_M) \Rightarrow 'c \Rightarrow 'a \text{ coG } \text{coF}$

where $\text{n2m_dtor_unfold_coG_coF } s1 \ s2 = \text{dtor_unfold_coF } (\text{BNF_Composition.id_bnf } o \ \text{BNF_Composition.id_bnf } o \ \text{map_pre_coGcoF}_M \ (\text{id} :: \text{unit} \Rightarrow \text{unit}) \ (\text{n2m_dtor_unfold_coG } s1 \ s2) \ \text{id } o \ s2)$

lemma $\text{coG_dtor_o_unfold} : \text{dtor_coG } o \ \text{dtor_unfold_coG } s = \text{map_pre_coG } \text{id } (\text{dtor_unfold_coG } s) \ o \ s$
<proof>

lemma $\text{coF_dtor_o_unfold} : \text{dtor_coF } o \ \text{dtor_unfold_coF } s = \text{map_pre_coF } \text{id } (\text{dtor_unfold_coF } s) \ o \ s$
<proof>

lemma $\text{coG_dtor_o_corec} : \text{dtor_coG } o \ \text{dtor_corec_coG } s = \text{map_pre_coG } \text{id } (\text{case_sum } \text{id } (\text{dtor_corec_coG } s)) \ o \ s$
<proof>

lemma $\text{coF_dtor_o_corec} : \text{dtor_coF } o \ \text{dtor_corec_coF } s = \text{map_pre_coF } \text{id } (\text{case_sum } \text{id } (\text{dtor_corec_coF } s)) \ o \ s$
<proof>

lemma $\text{n2m_dtor_unfold_coG} :$

$\text{dtor_coG } o \ \text{n2m_dtor_unfold_coG } s1 \ s2 = \text{BNF_Composition.id_bnf } o \ \text{BNF_Composition.id_bnf } o \ \text{map_pre_coG}_M \ \text{id } (\text{n2m_dtor_unfold_coG_coF } s1 \ s2) \ o \ s1$
<proof>

lemma $\text{n2m_dtor_unfold_coG_coF} :$

$\text{dtor_coF } o \ \text{n2m_dtor_unfold_coG_coF } s1 \ s2 = \text{BNF_Composition.id_bnf } o \ \text{BNF_Composition.id_bnf } o \ \text{map_pre_coGcoF}_M \ \text{id } (\text{n2m_dtor_unfold_coG } s1 \ s2) \ (\text{n2m_dtor_unfold_coG_coF } s1 \ s2) \ o \ s2$
<proof>

8.3.2 Corecursors

definition $\text{n2m_dtor_corec_coG} ::$

$('b \Rightarrow ('a \text{ coG } \text{coF} + 'c, 'a) \text{ coG}_M\text{_pre_coG}_M) \Rightarrow ('c \Rightarrow ('a \text{ coG } \text{coF} + 'c, 'a \text{ coG } + 'b) \text{ coGcoF}_M\text{_pre_coGcoF}_M) \Rightarrow 'b \Rightarrow 'a \text{ coG}$

where $\text{n2m_dtor_corec_coG } s1 \ s2 = \text{dtor_corec_coG } (\text{BNF_Composition.id_bnf } o \ \text{BNF_Composition.id_bnf } o \ \text{map_pre_coG}_M \ \text{id } (\text{id} :: \text{unit} \Rightarrow \text{unit}) \ (\text{case_sum } (\text{map_coF } \text{Inl}) \ (\text{dtor_corec_coF } (\text{BNF_Composition.id_bnf } o \ \text{BNF_Composition.id_bnf } o \ \text{map_pre_coGcoF}_M \ (\text{id} :: \text{unit} \Rightarrow \text{unit}) \ \text{id } (\text{map_sum } (\text{map_coF } \text{Inl}) \ \text{id}) \ o \ s2)))) \ o \ s1$

definition $\text{n2m_dtor_corec_coG_coF} ::$

$('b \Rightarrow ('a \text{ coG } \text{coF} + 'c, 'a) \text{ coG}_M\text{_pre_coG}_M) \Rightarrow ('c \Rightarrow ('a \text{ coG } \text{coF} + 'c, 'a \text{ coG } + 'b) \text{ coGcoF}_M\text{_pre_coGcoF}_M) \Rightarrow 'c \Rightarrow 'a \text{ coG } \text{coF}$

where $\text{n2m_dtor_corec_coG_coF } s1 \ s2 = \text{dtor_corec_coF } (\text{BNF_Composition.id_bnf } o \ \text{BNF_Composition.id_bnf } o \ \text{map_pre_coGcoF}_M \ (\text{id} :: \text{unit} \Rightarrow \text{unit}) \ (\text{case_sum } \text{id } (\text{n2m_dtor_corec_coG } s1 \ s2))) \ \text{id } o \ s2$

lemma $\text{n2m_dtor_corec_coG} :$

$\text{dtor_coG } o \ \text{n2m_dtor_corec_coG } s1 \ s2 = \text{BNF_Composition.id_bnf } o \ \text{BNF_Composition.id_bnf } o \ \text{map_pre_coG}_M \ \text{id } (\text{case_sum } \text{id } (\text{n2m_dtor_corec_coG_coF } s1 \ s2)) \ o \ s1$

<proof>

lemma *n2m_dtor_corec_coG_coF*:

dtor_coF o *n2m_dtor_corec_coG_coF* *s1 s2* = *BNF_Composition.id_bnf* o *BNF_Composition.id_bnf* o *map_pre_coGcoF_M*
id (*case_sum id* (*n2m_dtor_corec_coG* *s1 s2*)) (*case_sum id* (*n2m_dtor_corec_coG_coF* *s1 s2*)) o *s2*
<proof>

8.3.3 Coinduction

lemma *n2m_rel_coinduct_coG_coG_coF*:

assumes *CIH1*: $\forall x y. R x y \longrightarrow \text{BNF_Def.vimage2p } (\text{BNF_Composition.id_bnf} \circ \text{BNF_Composition.id_bnf})$
 $(\text{BNF_Composition.id_bnf} \circ \text{BNF_Composition.id_bnf}) (\text{rel_pre_coG}_M P R S) (\text{dtor_coG } x) (\text{dtor_coG } y)$

and *CIH2*: $\forall x y. S x y \longrightarrow \text{BNF_Def.vimage2p } (\text{BNF_Composition.id_bnf} \circ \text{BNF_Composition.id_bnf})$
 $(\text{BNF_Composition.id_bnf} \circ \text{BNF_Composition.id_bnf}) (\text{rel_pre_coGcoF}_M P R S) (\text{dtor_coF } x) (\text{dtor_coF } y)$

shows $R \leq \text{rel_coG } P \wedge S \leq \text{rel_coF } (\text{rel_coG } P)$

<proof>

lemmas *n2m_ctor_induct_coG_coG_coF* = *spec[OF spec[OF spec[OF spec[OF*

n2m_rel_coinduct_coG_coG_coF [*of* (=),

unfolded coG.rel_eq coF.rel_eq le_fun_def le_bool_def all_simps(1,2)[symmetric]]]]]]]