

Operations on Bounded Natural Functors

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Abstract

This entry formalizes the closure property of bounded natural functors (BNFs) under seven operations. These operations and the corresponding proofs constitute the core of Isabelle’s (co)datatype package. To be close to the implemented tactics, the proofs are deliberately formulated as detailed apply scripts. The (co)datatypes together with (co)induction principles and (co)recursors are byproducts of the fixpoint operations LFP and GFP. Composition of BNFs is subdivided into four simpler operations: Compose, Kill, Lift, and Permute. The N2M operation provides mutual (co)induction principles and (co)recursors for nested (co)datatypes.

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1 Least Fixpoint (a.k.a. Datatype)

ML *(open Ctr_Sugar_Util)*
notation *BNF_Def.convol* ($\langle _ , _ \rangle$)

'b1 = ('a, 'b1, 'b2) F1
'b2 = ('a, 'b1, 'b2) F2

To build a witness scenario, let us assume

('a, 'b1, 'b2) F1 = 'a * 'b1 + 'a * 'b2
('a, 'b1, 'b2) F2 = unit + 'b1 * 'b2

declare *[[bnf_internals]]*

bnf-axiomatization (*F1set1: 'a, F1set2: 'b1, F1set3: 'b2*) *F1*
[wits: 'a ⇒ 'b1 ⇒ ('a, 'b1, 'b2) F1 'a ⇒ 'b2 ⇒ ('a, 'b1, 'b2) F1]
for *map: F1map rel: F1rel*

bnf-axiomatization (*F2set1: 'a, F2set2: 'b1, F2set3: 'b2*) *F2*
[wits: ('a, 'b1, 'b2) F2]
for *map: F2map rel: F2rel*

abbreviation *F1in :: 'a1 set ⇒ 'a2 set ⇒ 'a3 set ⇒ (('a1, 'a2, 'a3) F1) set* **where**
F1in A1 A2 A3 ≡ {x. F1set1 x ⊆ A1 ∧ F1set2 x ⊆ A2 ∧ F1set3 x ⊆ A3}

abbreviation *F2in :: 'a1 set ⇒ 'a2 set ⇒ 'a3 set ⇒ (('a1, 'a2, 'a3) F2) set* **where**
F2in A1 A2 A3 ≡ {x. F2set1 x ⊆ A1 ∧ F2set2 x ⊆ A2 ∧ F2set3 x ⊆ A3}

lemma *F1map_comp_id: F1map g1 g2 g3 (F1map id f2 f3 x) = F1map g1 (g2 o f2) (g3 o f3) x*
apply *(rule trans)*
apply *(rule F1.map_comp)*
unfolding *o_id*
apply *(rule refl)*
done

lemmas *F1in_mono23 = F1.in_mono[OF subset_refl]*

lemma *F1map_congL: [∀ a ∈ F1set2 x. f a = a; ∀ a ∈ F1set3 x. g a = a] ⇒⇒*
F1map id f g x = x
apply *(rule trans)*
apply *(rule F1.map_cong0)*
apply *(rule refl)*
apply *(rule trans)*
apply *(erule bspec)*
apply *assumption*
apply *(rule sym)*
apply *(rule id_apply)*
apply *(rule trans)*
apply *(erule bspec)*
apply *assumption*
apply *(rule sym)*
apply *(rule id_apply)*
apply *(rule F1.map_id)*
done

lemma *F2map_comp_id: F2map g1 g2 g3 (F2map id f2 f3 x) = F2map g1 (g2 o f2) (g3 o f3) x*

```

apply (rule trans)
apply (rule F2.map_comp)
unfolding o_id
apply (rule refl)
done

```

lemmas F2in_mono23 = F2.in_mono[OF subset_refl]

```

lemma F2map_congL:  $[\forall a \in F2set2\ x. f\ a = a; \forall a \in F2set3\ x. g\ a = a] \implies$ 
  F2map id f g x = x
apply (rule trans)
apply (rule F2.map_cong0)
apply (rule refl)
apply (rule trans)
apply (erule bspec)
apply assumption
apply (rule sym)
apply (rule id_apply)
apply (rule trans)
apply (erule bspec)
apply assumption
apply (rule sym)
apply (rule id_apply)
apply (rule F2.map_id)
done

```

1.1 Algebra

definition alg where

```

alg B1 B2 s1 s2 =
  (( $\forall x \in F1in\ (UNIV :: 'a\ set)\ B1\ B2. s1\ x \in B1$ )  $\wedge$  ( $\forall y \in F2in\ (UNIV :: 'a\ set)\ B1\ B2. s2\ y \in B2$ ))

```

```

lemma alg_F1set:  $[\text{alg } B1\ B2\ s1\ s2; F1set2\ x \subseteq B1; F1set3\ x \subseteq B2] \implies s1\ x \in B1$ 
apply (tactic  $\langle\langle$  dtac @{context} @{thm iffD1[OF alg_def]} 1  $\rangle\rangle$ )
apply (erule conjE)+
apply (erule bspec)
apply (rule CollectI)
apply (rule conjI[OF subset_UNIV])
apply (erule conjI)
apply assumption
done

```

```

lemma alg_F2set:  $[\text{alg } B1\ B2\ s1\ s2; F2set2\ x \subseteq B1; F2set3\ x \subseteq B2] \implies s2\ x \in B2$ 
apply (tactic  $\langle\langle$  dtac @{context} @{thm iffD1[OF alg_def]} 1  $\rangle\rangle$ )
apply (erule conjE)+
apply (erule bspec)
apply (rule CollectI)
apply (rule conjI[OF subset_UNIV])
apply (erule conjI)
apply assumption
done

```

```

lemma alg_not_empty:
  alg B1 B2 s1 s2  $\implies B1 \neq \{\} \wedge B2 \neq \{\}$ 
apply (rule conjI)
apply (rule notI)
apply (tactic  $\langle\langle$  hyp_subst_tac @{context} 1  $\rangle\rangle$ )
apply (erule alg_F1set)

```

```

apply (rule subset_emptyI)
apply (erule F1.wit1 F1.wit2 F2.wit)

```

apply (rule subsetI)
apply (drule F1.wit1 F1.wit2 F2.wit)

apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (tactic << FIRST' (map (fn thm => rtac @ {context} thm THEN' assume_tac @ {context}) @ {thms
alg_F1set alg_F2set}) 1 >>))

apply (rule subset_emptyI)
apply (erule F1.wit1 F1.wit2 F2.wit)

apply (rule subsetI)
apply (drule F1.wit1 F1.wit2 F2.wit)
apply (erule FalseE)

apply (erule emptyE)

apply (rule notI)
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (drule alg_F2set)

apply (rule subsetI)
apply (rule FalseE)
apply (erule F1.wit1 F1.wit2 F2.wit)

apply (rule subset_emptyI)
apply (erule F1.wit1 F1.wit2 F2.wit)

apply (erule emptyE)
done

1.2 Morphism

definition mor where

$$\begin{aligned} \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g = \\ & ((\forall a \in B1. f \ a \in B1') \wedge (\forall a \in B2. g \ a \in B2')) \wedge \\ & ((\forall z \in F1in \ (UNIV :: 'a \ set) \ B1 \ B2. f \ (s1 \ z) = s1' \ (F1map \ id \ f \ g \ z)) \wedge \\ & (\forall z \in F2in \ (UNIV :: 'a \ set) \ B1 \ B2. g \ (s2 \ z) = s2' \ (F2map \ id \ f \ g \ z))) \end{aligned}$$

lemma morE1: $\llbracket \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g; z \in F1in \ UNIV \ B1 \ B2 \rrbracket$
 $\implies f \ (s1 \ z) = s1' \ (F1map \ id \ f \ g \ z)$
apply (tactic << dtac @ {context} @ {thm iffD1[OF mor_def]} 1 >>)
apply (erule conjE)+
apply (erule bspec)
apply assumption
done

lemma morE2: $\llbracket \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g; z \in F2in \ UNIV \ B1 \ B2 \rrbracket$
 $\implies g \ (s2 \ z) = s2' \ (F2map \ id \ f \ g \ z)$
apply (tactic << dtac @ {context} @ {thm iffD1[OF mor_def]} 1 >>)
apply (erule conjE)+
apply (erule bspec)
apply assumption
done

lemma mor_incl: $\llbracket B1 \subseteq B1'; B2 \subseteq B2' \rrbracket \implies \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1 \ s2 \ id \ id$
apply (tactic << rtac @ {context} (@ {thm mor_def} RS iffD2) 1 >>)
apply (rule conjI)

apply (rule conjI)
apply (rule ballI)
apply (erule set_mp)
apply (erule ssubst_mem[OF id_apply])

```

apply (rule ballI)
apply (erule set_mp)
apply (erule ssubst_mem[OF id_apply])

```

```

apply (rule conjI)
apply (rule ballI)
apply (rule trans)
  apply (rule id_apply)
apply (tactic << stac @ {context} @ {thm F1.map_id} 1 >>)
apply (rule refl)

```

```

apply (rule ballI)
apply (rule trans)
  apply (rule id_apply)
apply (tactic << stac @ {context} @ {thm F2.map_id} 1 >>)
apply (rule refl)
done

```

lemma mor_comp:

```

[[mor B1 B2 s1 s2 B1' B2' s1' s2' f g;
  mor B1' B2' s1' s2' B1'' B2'' s1'' s2'' f' g]] ==>
mor B1 B2 s1 s2 B1'' B2'' s1'' s2'' (f' o f) (g' o g)
apply (tactic << dtac @ {context} (@ {thm mor_def} RS iffD1) 1 >>)
apply (tactic << dtac @ {context} (@ {thm mor_def} RS iffD1) 1 >>)
apply (tactic << rtac @ {context} (@ {thm mor_def} RS iffD2) 1 >>)
apply (erule conjE)+
apply (rule conjI)

```

```

apply (rule conjI)
apply (rule ballI)
apply (rule ssubst_mem[OF o_apply])
apply (erule bspec)
apply (erule bspec)
apply assumption

```

```

apply (rule ballI)
apply (rule ssubst_mem[OF o_apply])
apply (erule bspec)
apply (erule bspec)
apply assumption

```

```

apply (rule conjI)
apply (rule ballI)
apply (rule trans[OF o_apply])
apply (rule trans)
apply (rule trans)
  apply (erule bspec[rotated])
  apply assumption
  apply (erule arg_cong)
apply (erule CollectE conjE)+
apply (erule bspec)
apply (rule CollectI)
apply (rule conjI)
  apply (rule subset_UNIV)
apply (rule conjI)
  apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(2))
  apply (rule image_subsetI)
  apply (erule bspec)
  apply (erule set_mp)
  apply assumption
apply (rule ord_eq_le_trans)

```

```

  apply (rule F1.set_map(3))
  apply (rule image_subsetI)
  apply (erule bspec)
  apply (erule set_mp)
  apply assumption
  apply (rule arg_cong[OF F1map_comp_id])

```

```

apply (rule ballI)
apply (rule trans[OF o_apply])
apply (rule trans)
apply (rule trans)
  apply (drule bspec[rotated])
  apply assumption
  apply (erule arg_cong)
  apply (erule CollectE conjE)+
  apply (erule bspec)
  apply (rule CollectI)
  apply (rule conjI)
  apply (rule subset_UNIV)
  apply (rule conjI)
  apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(2))
  apply (rule image_subsetI)
  apply (erule bspec)
  apply (erule set_mp)
  apply assumption
  apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(3))
  apply (rule image_subsetI)
  apply (erule bspec)
  apply (erule set_mp)
  apply assumption
  apply (rule arg_cong[OF F2map_comp_id])
done

```

lemma *mor_cong*: $\llbracket f' = f; g' = g; \text{mor } B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f\ g \rrbracket \implies \text{mor } B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ f'\ g'$

```

  apply (tactic  $\ll \text{hyp\_subst\_tac } @\{\text{context}\} 1 \gg$ )
  apply assumption
done

```

lemma *mor_str*:

```

  mor UNIV UNIV (F1map id s1 s2) (F2map id s1 s2) UNIV UNIV s1 s2 s1 s2
  apply (rule iffD2)
  apply (rule mor_def)
  apply (rule conjI)
  apply (rule conjI)
  apply (rule ballI)
  apply (rule UNIV_I)
  apply (rule ballI)
  apply (rule UNIV_I)

```

```

  apply (rule conjI)
  apply (rule ballI)
  apply (rule refl)
  apply (rule ballI)
  apply (rule refl)
done

```

1.3 Bounds

```

type-synonym bd_type_F1' = bd_type_F1 + (bd_type_F1, bd_type_F1, bd_type_F1) F1
type-synonym bd_type_F2' = bd_type_F2 + (bd_type_F2, bd_type_F2, bd_type_F2) F2
type-synonym SucFbd_type = ((bd_type_F1' + bd_type_F2') set)

```

type-synonym 'a1 ASucFbd_type = (SucFbd_type \Rightarrow ('a1 + bool))

abbreviation F1bd' \equiv bd_F1 +c |UNIV :: (bd_type_F1, bd_type_F1, bd_type_F1) F1 set|

lemma F1set1_bd_incr: $\bigwedge x. |F1set1\ x| \leq_o F1bd'$

by (rule ordLeq_transitive[OF F1.set_bd(1) ordLeq_csum1[OF F1.bd_Card_order]])

lemma F1set2_bd_incr: $\bigwedge x. |F1set2\ x| \leq_o F1bd'$

by (rule ordLeq_transitive[OF F1.set_bd(2) ordLeq_csum1[OF F1.bd_Card_order]])

lemma F1set3_bd_incr: $\bigwedge x. |F1set3\ x| \leq_o F1bd'$

by (rule ordLeq_transitive[OF F1.set_bd(3) ordLeq_csum1[OF F1.bd_Card_order]])

lemmas F1bd'_Card_order = Card_order_csum

lemmas F1bd'_Cinfinite = Cinfinite_csum1[OF F1.bd_Cinfinite]

lemmas F1bd'_Cnotzero = Cinfinite_Cnotzero[OF F1bd'_Cinfinite]

lemmas F1bd'_card_order = card_order_csum[OF F1.bd_card_order card_of_card_order_on]

abbreviation F2bd' \equiv bd_F2 +c |UNIV :: (bd_type_F2, bd_type_F2, bd_type_F2) F2 set|

lemma F2set1_bd_incr: $\bigwedge x. |F2set1\ x| \leq_o F2bd'$

by (rule ordLeq_transitive[OF F2.set_bd(1) ordLeq_csum1[OF F2.bd_Card_order]])

lemma F2set2_bd_incr: $\bigwedge x. |F2set2\ x| \leq_o F2bd'$

by (rule ordLeq_transitive[OF F2.set_bd(2) ordLeq_csum1[OF F2.bd_Card_order]])

lemma F2set3_bd_incr: $\bigwedge x. |F2set3\ x| \leq_o F2bd'$

by (rule ordLeq_transitive[OF F2.set_bd(3) ordLeq_csum1[OF F2.bd_Card_order]])

lemmas F2bd'_Card_order = Card_order_csum

lemmas F2bd'_Cinfinite = Cinfinite_csum1[OF F2.bd_Cinfinite]

lemmas F2bd'_Cnotzero = Cinfinite_Cnotzero[OF F2bd'_Cinfinite]

lemmas F2bd'_card_order = card_order_csum[OF F2.bd_card_order card_of_card_order_on]

abbreviation SucFbd **where** SucFbd \equiv cardSuc (F1bd' +c F2bd')

abbreviation ASucFbd **where** ASucFbd \equiv (|UNIV| +c ctwo) \wedge c SucFbd

lemma F1set1_bd': $|F1set1\ x| \leq_o F1bd' +c F2bd'$

apply (rule ordLeq_transitive)

apply (rule F1set1_bd_incr)

apply (rule ordLeq_csum1)

apply (rule F1bd'_Card_order)

done

lemma F1set2_bd': $|F1set2\ x| \leq_o F1bd' +c F2bd'$

apply (rule ordLeq_transitive)

apply (rule F1set2_bd_incr)

apply (rule ordLeq_csum1)

apply (rule F1bd'_Card_order)

done

lemma F1set3_bd': $|F1set3\ x| \leq_o F1bd' +c F2bd'$

apply (rule ordLeq_transitive)

apply (rule F1set3_bd_incr)

apply (rule ordLeq_csum1)

apply (rule F1bd'_Card_order)

done

lemma F2set1_bd': $|F2set1\ x| \leq_o F1bd' +c F2bd'$

apply (rule ordLeq_transitive)

apply (rule F2set1_bd_incr)

apply (rule ordLeq_csum2)

apply (rule F2bd'_Card_order)

done

lemma F2set2_bd': $|F2set2\ x| \leq_o F1bd' +c F2bd'$

apply (rule ordLeq_transitive)

apply (rule F2set2_bd_incr)

apply (rule ordLeq_csum2)

apply (rule *F2bd'_Card_order*)
done

lemma *F2set3_bd'*: $|F2set3\ x| \leq_o F1bd' + c\ F2bd'$
apply (rule *ordLeq_transitive*)
apply (rule *F2set3_bd_incr*)
apply (rule *ordLeq_csum2*)
apply (rule *F2bd'_Card_order*)
done

lemmas *SucFbd_Card_order* = *cardSuc_Card_order*[*OF Card_order_csum*]
lemmas *SucFbd_Cinfinite* = *Cinfinite_cardSuc*[*OF Cinfinite_csum1*[*OF F1bd'_Cinfinite*]]
lemmas *SucFbd_Cnotzero* = *Cinfinite_Cnotzero*[*OF SucFbd_Cinfinite*]
lemmas *worel_SucFbd* = *Card_order_wo_rel*[*OF SucFbd_Card_order*]
lemmas *ASucFbd_Cinfinite* = *Cinfinite_cexp*[*OF ordLeq_csum2*[*OF Card_order_ctwo*] *SucFbd_Cinfinite*]

1.4 Minimal Algebras

abbreviation *min_G1* **where**
 $min_G1\ As1_As2\ i \equiv (\bigcup j \in underS\ SucFbd\ i.\ fst\ (As1_As2\ j))$

abbreviation *min_G2* **where**
 $min_G2\ As1_As2\ i \equiv (\bigcup j \in underS\ SucFbd\ i.\ snd\ (As1_As2\ j))$

abbreviation *min_H* **where**
 $min_H\ s1\ s2\ As1_As2\ i \equiv$
 $(min_G1\ As1_As2\ i \cup s1\ ' (F1in\ (UNIV\ ::\ 'a\ set)\ (min_G1\ As1_As2\ i)\ (min_G2\ As1_As2\ i)),$
 $min_G2\ As1_As2\ i \cup s2\ ' (F2in\ (UNIV\ ::\ 'a\ set)\ (min_G1\ As1_As2\ i)\ (min_G2\ As1_As2\ i)))$

abbreviation *min_algs* **where**
 $min_algs\ s1\ s2 \equiv wo_rel.worec\ SucFbd\ (min_H\ s1\ s2)$

definition *min_alg1* **where**
 $min_alg1\ s1\ s2 = (\bigcup i \in Field\ SucFbd.\ fst\ (min_algs\ s1\ s2\ i))$

definition *min_alg2* **where**
 $min_alg2\ s1\ s2 = (\bigcup i \in Field\ SucFbd.\ snd\ (min_algs\ s1\ s2\ i))$

lemma *min_algs*:
 $i \in Field\ SucFbd \implies min_algs\ s1\ s2\ i = min_H\ s1\ s2\ (min_algs\ s1\ s2)\ i$
apply (rule *fun_cong*[*OF wo_rel.worec_fixpoint*[*OF worel_SucFbd*]])
apply (rule *iffD2*)
apply (rule *meta_eq_to_obj_eq*)
apply (rule *wo_rel.adm_wo_def*[*OF worel_SucFbd*])
apply (rule *allI*)
apply (rule *impI*)

apply (rule *iffD2*)
apply (rule *prod.inject*)
apply (rule *conjI*)

apply (rule *arg_cong2*[*of _ _ _ _ (U)*])
apply (rule *SUP_cong*)
apply (rule *refl*)
apply (rule *bspec*)
apply (rule *assumption*)
apply (rule *arg_cong*)

apply (rule *image_cong*)
apply (rule *arg_cong2*[*of _ _ _ _ F1in UNIV*])
apply (rule *SUP_cong*)
apply (rule *refl*)
apply (rule *bspec*)
apply (rule *assumption*)


```

    apply (erule arg_cong)
  apply (rule SUP_cong)
  apply (rule refl)
  apply (drule bspec)
  apply assumption
  apply (erule arg_cong)
  apply (rule refl)

  apply (rule arg_cong2[of _ _ _ _ (∪)])
  apply (rule SUP_cong)
  apply (rule refl)
  apply (drule bspec)
  apply assumption
  apply (erule arg_cong)

  apply (rule image_cong)
  apply (rule arg_cong2[of _ _ _ _ F2in UNIV])
  apply (rule SUP_cong)
  apply (rule refl)
  apply (drule bspec)
  apply assumption
  apply (erule arg_cong)
  apply (rule SUP_cong)
  apply (rule refl)
  apply (drule bspec)
  apply assumption
  apply (erule arg_cong)
  apply (rule refl)
done

corollary min_algs1: i ∈ Field SucFbd ⇒ fst (min_algs s1 s2 i) =
  min_G1 (min_algs s1 s2) i ∪
  s1 ‘ (F1in UNIV (min_G1 (min_algs s1 s2) i) (min_G2 (min_algs s1 s2) i))
  apply (rule trans)
  apply (erule arg_cong[OF min_algs])
  apply (rule fst_conv)
done

corollary min_algs2: i ∈ Field SucFbd ⇒ snd (min_algs s1 s2 i) =
  min_G2 (min_algs s1 s2) i ∪
  s2 ‘ (F2in UNIV (min_G1 (min_algs s1 s2) i) (min_G2 (min_algs s1 s2) i))
  apply (rule trans)
  apply (erule arg_cong[OF min_algs])
  apply (rule snd_conv)
done

lemma min_algs_mono1: relChain SucFbd (%i. fst (min_algs s1 s2 i))
  apply (tactic ⟨⟨ rtac @ {context} @ {thm iffD2[OF meta_eq_to_obj_eq[OF relChain_def]]} 1 ⟩⟩)
  apply (rule allI)+
  apply (rule impI)
  apply (rule case_split)
  apply (rule xt1(3))
  apply (rule min_algs1)
  apply (erule FieldI2)
  apply (rule subsetI)
  apply (rule UnI1)
  apply (rule UN_I)
  apply (erule underS_I)
  apply assumption
  apply assumption
  apply (rule equalityD1)
  apply (drule notnotD)
  apply (erule arg_cong)

```

done

```
lemma min_algs_mono2: relChain SucFbd (%i. snd (min_algs s1 s2 i))
  apply (tactic << rtac @{context} @{thm iffD2[OF meta_eq_to_obj_eq[OF relChain_def]] 1 >>)
  apply (rule allI)+
  apply (rule impI)
  apply (rule case_split)
  apply (rule xt1(3))
  apply (rule min_algs2)
  apply (erule FieldI2)
  apply (rule subsetI)
  apply (rule UnI1)
  apply (rule UN_I)
  apply (erule underS_I)
  apply assumption
  apply assumption
  apply (rule equalityD1)
  apply (drule notnotD)
  apply (erule arg_cong)
done
```

```
lemma SucFbd_limit: [x1 ∈ Field SucFbd & x2 ∈ Field SucFbd]
  ⇒ ∃ y ∈ Field SucFbd. (x1 ≠ y ∧ (x1, y) ∈ SucFbd) ∧ (x2 ≠ y ∧ (x2, y) ∈ SucFbd)
  apply (erule conjE)+
  apply (rule rev_mp)
  apply (rule Cinfinites_limit_finite)
  apply (rule finite.insertI)
  apply (rule finite.insertI)
  apply (rule finite.emptyI)
  apply (erule insert_subsetI)
  apply (erule insert_subsetI)
  apply (rule empty_subsetI)
  apply (rule SucFbd_Cinfinites)
  apply (rule impI)
  apply (erule bexE)
  apply (rule bexE)

  apply (rule conjI)

  apply (erule bspec)
  apply (rule insertI1)

  apply (erule bspec)
  apply (rule insertI2)
  apply (rule insertI1)
  apply assumption
done
```

```
lemma alg_min_alg: alg (min_alg1 s1 s2) (min_alg2 s1 s2) s1 s2
  apply (tactic << rtac @{context} (@{thm alg_def} RS iffD2) 1 >>)
  apply (rule conjI)
  apply (rule ballI)
  apply (erule CollectE conjE)+

  apply (rule bexE)
  apply (rule cardSuc_UNION_Cinfinites)
  apply (rule Cinfinites_csum1)
  apply (rule F1bd'_Cinfinites)
  apply (rule min_algs_mono1)
  apply (erule subset_trans[OF _ equalityD1[OF min_alg1_def]])
  apply (rule ordLeq_transitive)
  apply (rule F1set2_bd_incr)
  apply (rule ordLeq_csum1)
```

```

apply (rule F1bd'_ Card_order)

apply (rule bexE)
apply (rule cardSuc_UNION_Cinfinite)
  apply (rule Cinfinite_csum1)
  apply (rule F1bd'_ Cinfinite)
  apply (rule min_algs_mono2)
apply (erule subset_trans[OF equalityD1[OF min_alg2_def]])
apply (rule ordLeq_transitive)
  apply (rule F1set3_bd_incr)
apply (rule ordLeq_csum1)
apply (rule F1bd'_ Card_order)

apply (rule bexE)
apply (rule SucFbd_limit)
apply (erule conjI)
apply assumption
apply (rule set_mp[OF equalityD2[OF min_alg1_def]])
apply (rule UN_I)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply assumption
apply (rule set_mp)
apply (rule equalityD2)
apply (rule min_algs1)
apply assumption
apply (rule UnI2)
apply (rule image_eqI)
apply (rule refl)
apply (rule CollectI)
apply (erule asm_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule conjE)+

apply (rule conjI)
apply assumption

apply (rule conjI)
apply (erule subset_trans)
apply (rule subsetI)
apply (rule UN_I)
  apply (erule underS_I)
  apply assumption
apply assumption

apply (erule subset_trans)
apply (erule UN_upper[OF underS_I])
apply assumption

apply (rule ballI)
apply (erule CollectE conjE)+

apply (rule bexE)
apply (rule cardSuc_UNION_Cinfinite)
  apply (rule Cinfinite_csum1)

```

```

apply (rule F1bd'_Cinfinite)
apply (rule min_algs_mono1)

apply (erule subset_trans[OF _ equalityD1[OF min_alg1_def]])
apply (rule ordLeq_transitive)
apply (rule F2set2_bd_incr)
apply (rule ordLeq_csum2)
apply (rule F2bd'_Card_order)

apply (rule bexE)
apply (rule cardSuc_UNION_Cinfinite)
apply (rule Cinfinite_csum1)
apply (rule F1bd'_Cinfinite)
apply (rule min_algs_mono2)

apply (erule subset_trans[OF _ equalityD1[OF min_alg2_def]])
apply (rule ordLeq_transitive)
apply (rule F2set3_bd_incr)
apply (rule ordLeq_csum2)
apply (rule F2bd'_Card_order)

apply (rule bexE)
apply (rule SucFbd_limit)
apply (erule conjI)
apply assumption
apply (rule set_mp[OF equalityD2[OF min_alg2_def]])
apply (rule UN_I)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply assumption
apply (rule set_mp)
apply (rule equalityD2)
apply (rule min_algs2)
apply assumption
apply (rule UnI2)
apply (rule image_eqI)
apply (rule refl)
apply (rule CollectI)
apply (rule conjI)
apply assumption

apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule conjE)+
apply (rule conjI)
apply (erule subset_trans)
apply (rule UN_upper)
apply (erule underS_I)
apply assumption

apply (erule subset_trans)
apply (rule UN_upper)
apply (erule underS_I)
apply assumption
done

```

lemmas SucFbd_ASucFbd = ordLess_ordLeq_trans[OF

ordLess_ctwo_cexp
cexp_mono1[*OF ordLeq_csum2*[*OF Card_order_ctwo*]],
OF SucFbd_Card_order SucFbd_Card_order]

lemma *card_of_min_algs*:
fixes *s1* :: ('a, 'b, 'c) *F1* \Rightarrow 'b **and** *s2* :: ('a, 'b, 'c) *F2* \Rightarrow 'c
shows *i* \in *Field SucFbd* \longrightarrow
(|fst (*min_algs s1 s2 i*)| \leq_o (*ASucFbd* :: 'a *ASucFbd_type rel*) \wedge |snd (*min_algs s1 s2 i*)| \leq_o (*ASucFbd* :: 'a *ASucFbd_type rel*))
apply (*rule well_order_induct_imp*[*of* _ %*i*. (|fst (*min_algs s1 s2 i*)| \leq_o *ASucFbd* \wedge |snd (*min_algs s1 s2 i*)| \leq_o *ASucFbd*), *OF wored_SucFbd*])
apply (*rule impI*)
apply (*rule conjI*)
apply (*rule ordIso_ordLeq_trans*)
apply (*rule card_of_ordIso_subst*)
apply (*erule min_algs1*)
apply (*rule Un_Cinfinite_bound*)

apply (*rule UNION_Cinfinite_bound*)

apply (*rule ordLess_imp_ordLeq*)
apply (*rule ordLess_transitive*)
apply (*rule card_of_underS*)
apply (*rule SucFbd_Card_order*)
apply *assumption*
apply (*rule SucFbd_ASucFbd*)

apply (*rule ballI*)
apply (*erule allE*)
apply (*drule mp*)
apply (*erule underS_E*)
apply (*drule mp*)
apply (*erule underS_Field*)
apply (*erule conjE*)
apply *assumption*

apply (*rule ASucFbd_Cinfinite*)

apply (*rule ordLeq_transitive*)
apply (*rule card_of_image*)
apply (*rule ordLeq_transitive*)
apply (*rule F1.in_bd*)
apply (*rule ordLeq_transitive*)
apply (*rule cexp_mono1*)
apply (*rule csum_mono1*)
apply (*rule csum_mono2*)
apply (*rule csum_cinfinite_bound*)
apply (*rule UNION_Cinfinite_bound*)

apply (*rule ordLess_imp_ordLeq*)
apply (*rule ordLess_transitive*)
apply (*rule card_of_underS*)
apply (*rule SucFbd_Card_order*)
apply *assumption*
apply (*rule SucFbd_ASucFbd*)

apply (*rule ballI*)
apply (*erule allE*)
apply (*drule mp*)
apply (*erule underS_E*)
apply (*drule mp*)
apply (*erule underS_Field*)
apply (*erule conjE*)

```

apply assumption

apply (rule ASucFbd_Cinfinite)

apply (rule UNION_Cinfinite_bound)

apply (rule ordLess_imp_ordLeq)
apply (rule ordLess_transitive)
apply (rule card_of_underS)
apply (rule SucFbd_Card_order)
apply assumption
apply (rule SucFbd_ASucFbd)

apply (rule ballI)
apply (erule allE)
apply (drule mp)
apply (erule underS_E)
apply (drule mp)
apply (erule underS_Field)
apply (erule conjE)+
apply assumption

apply (rule ASucFbd_Cinfinite)

apply (rule card_of_Card_order)
apply (rule card_of_Card_order)
apply (rule ASucFbd_Cinfinite)

apply (rule F1bd'_Card_order)
apply (rule ordIso_ordLeq_trans)
apply (rule cexp_cong1)
apply (rule ordIso_transitive)
apply (rule csum_cong1)
apply (rule ordIso_transitive)
apply (tactic ⟨⟨ BNF_Tactics.mk_rotate_eq_tac @{context}
  (rtac @{context} @{thm ordIso_refl}) THEN'
  FIRST' [rtac @{context} @{thm card_of_Card_order}],
  rtac @{context} @{thm Card_order_csum},
  rtac @{context} @{thm Card_order_cexp}]]
  @{thm ordIso_transitive} @{thm csum_assoc} @{thm csum_com} @{thm csum_cong}
  [1,2] [2,1] 1 ⟩⟩)
apply (rule csum_absorb1)
apply (rule ASucFbd_Cinfinite)

apply (rule ordLeq_transitive)
apply (rule ordLeq_csum1)
apply (tactic ⟨⟨ FIRST' [rtac @{context} @{thm Card_order_csum}], rtac @{context} @{thm card_of_Card_order}]
1 ⟩⟩)
apply (rule ordLeq_cexp1)
apply (rule SucFbd_Cnotzero)
apply (rule Card_order_csum)
apply (rule csum_absorb1)
apply (rule ASucFbd_Cinfinite)
apply (rule ctwo_ordLeq_Cinfinite)
apply (rule ASucFbd_Cinfinite)
apply (rule F1bd'_Card_order)
apply (rule ordIso_imp_ordLeq)
apply (rule cexp_cprod_ordLeq)

apply (rule Card_order_csum)
apply (rule SucFbd_Cinfinite)
apply (rule F1bd'_Cnotzero)
apply (rule ordLeq_transitive)

```

```

apply (rule ordLeq_csum1)
apply (rule F1bd'_Card_order)
apply (rule cardSuc_ordLeq)
apply (rule Card_order_csum)

apply (rule ASucFbd_Cinfinite)

apply (rule ordIso_ordLeq_trans)
apply (rule card_of_ordIso_subst)
apply (erule min_algs2)
apply (rule Un_Cinfinite_bound)

apply (rule UNION_Cinfinite_bound)

apply (rule ordLess_imp_ordLeq)
apply (rule ordLess_transitive)
apply (rule card_of_underS)
apply (rule SucFbd_Card_order)
apply assumption
apply (rule SucFbd_ASucFbd)

apply (rule ballI)
apply (erule allE)
apply (drule mp)
apply (erule underS_E)
apply (drule mp)
apply (erule underS_Field)
apply (erule conjE)+
apply assumption

apply (rule ASucFbd_Cinfinite)

apply (rule ordLeq_transitive)
apply (rule card_of_image)
apply (rule ordLeq_transitive)
apply (rule F2.in_bd)
apply (rule ordLeq_transitive)
apply (rule cexp_mono1)
apply (rule csum_mono1)
apply (rule csum_mono2)
apply (rule csum_cinfinite_bound)
apply (rule UNION_Cinfinite_bound)

apply (rule ordLess_imp_ordLeq)
apply (rule ordLess_transitive)
apply (rule card_of_underS)
apply (rule SucFbd_Card_order)
apply assumption
apply (rule SucFbd_ASucFbd)

apply (rule ballI)
apply (erule allE)
apply (drule mp)
apply (erule underS_E)
apply (drule mp)
apply (erule underS_Field)
apply (erule conjE)+
apply assumption

apply (rule ASucFbd_Cinfinite)

apply (rule UNION_Cinfinite_bound)

```

```

apply (rule ordLess_imp_ordLeq)
apply (rule ordLess_transitive)
apply (rule card_of_underS)
  apply (rule SucFbd_Card_order)
apply assumption
apply (rule SucFbd_ASucFbd)

apply (rule ballI)
apply (erule allE)
apply (drule mp)
  apply (erule underS_E)
apply (drule mp)
  apply (erule underS_Field)
apply (erule conjE)+
apply assumption

apply (rule ASucFbd_Cinfinite)

apply (rule card_of_Card_order)
apply (rule card_of_Card_order)
apply (rule ASucFbd_Cinfinite)

apply (rule F2bd'_Card_order)
apply (rule ordIso_ordLeq_trans)
apply (rule cexp_cong1)

apply (rule ordIso_transitive)
apply (rule csum_cong1)
apply (rule ordIso_transitive)
apply (tactic ⟨⟨ BNF_Tactics.mk_rotate_eq_tac @ {context}
  (rtac @ {context} @ {thm ordIso_refl} THEN'
  FIRST' [rtac @ {context} @ {thm card_of_Card_order},
  rtac @ {context} @ {thm Card_order_csum},
  rtac @ {context} @ {thm Card_order_cexp}])
  @ {thm ordIso_transitive} @ {thm csum_assoc} @ {thm csum_com} @ {thm csum_cong}
  [1,2] [2,1] 1 ⟩⟩)
apply (rule csum_absorb1)
apply (rule ASucFbd_Cinfinite)

apply (rule ordLeq_transitive)
apply (rule ordLeq_csum1)
apply (tactic ⟨⟨ FIRST' [rtac @ {context} @ {thm Card_order_csum}, rtac @ {context} @ {thm card_of_Card_order}]
1 ⟩⟩)
apply (rule ordLeq_cexp1)
apply (rule SucFbd_Cnotzero)
apply (rule Card_order_csum)

apply (rule csum_absorb1)
apply (rule ASucFbd_Cinfinite)
apply (rule ctwo_ordLeq_Cinfinite)
apply (rule ASucFbd_Cinfinite)
apply (rule F2bd'_Card_order)
apply (rule ordIso_imp_ordLeq)
apply (rule cexp_cprod_ordLeq)
  apply (rule Card_order_csum)
  apply (rule SucFbd_Cinfinite)
apply (rule F2bd'_Cnotzero)
apply (rule ordLeq_transitive)
apply (rule ordLeq_csum2)
apply (rule F2bd'_Card_order)
apply (rule cardSuc_ordLeq)
apply (rule Card_order_csum)

```



```

apply (rule ASucFbd_Cinfinite)
done

```

```

lemma card_of_min_alg1:
fixes s1 :: ('a, 'b, 'c) F1  $\Rightarrow$  'b and s2 :: ('a, 'b, 'c) F2  $\Rightarrow$  'c
shows |min_alg1 s1 s2|  $\leq$  o (ASucFbd :: 'a ASucFbd_type rel)
apply (rule ordIso_ordLeq_trans)
apply (rule card_of_ordIso_subst[OF min_alg1_def])
apply (rule UNION_Cinfinite_bound)

```

```

apply (rule ordIso_ordLeq_trans)
apply (rule card_of_Field_ordIso)
apply (rule SucFbd_Card_order)
apply (rule ordLess_imp_ordLeq)
apply (rule SucFbd_ASucFbd)

```

```

apply (rule ballI)
apply (drule rev_mp)
apply (rule card_of_min_algs)
apply (erule conjE)+
apply assumption
apply (rule ASucFbd_Cinfinite)
done

```

```

lemma card_of_min_alg2:
fixes s1 :: ('a, 'b, 'c) F1  $\Rightarrow$  'b and s2 :: ('a, 'b, 'c) F2  $\Rightarrow$  'c
shows |min_alg2 s1 s2|  $\leq$  o (ASucFbd :: 'a ASucFbd_type rel)
apply (rule ordIso_ordLeq_trans)
apply (rule card_of_ordIso_subst[OF min_alg2_def])
apply (rule UNION_Cinfinite_bound)

```

```

apply (rule ordIso_ordLeq_trans)
apply (rule card_of_Field_ordIso)
apply (rule SucFbd_Card_order)
apply (rule ordLess_imp_ordLeq)
apply (rule SucFbd_ASucFbd)

```

```

apply (rule ballI)
apply (drule rev_mp)
apply (rule card_of_min_algs)
apply (erule conjE)+
apply assumption
apply (rule ASucFbd_Cinfinite)
done

```

```

lemma least_min_algs: alg B1 B2 s1 s2  $\implies$ 
  i  $\in$  Field SucFbd  $\longrightarrow$ 

```

```

  fst (min_algs s1 s2 i)  $\subseteq$  B1  $\wedge$  snd (min_algs s1 s2 i)  $\subseteq$  B2

```

```

apply (rule well_order_induct_imp[of _ %i. (fst (min_algs s1 s2 i)  $\subseteq$  B1  $\wedge$  snd (min_algs s1 s2 i)  $\subseteq$  B2), OF
worel_SucFbd])

```

```

apply (rule impI)
apply (rule conjI)
apply (rule ord_eq_le_trans)
apply (erule min_algs1)
apply (rule Un_least)
apply (rule UN_least)
apply (erule allE)
apply (drule mp)
apply (erule underS_E)
apply (drule mp)
apply (erule underS_Field)
apply (erule conjE)+
apply assumption

```

```

apply (rule image_subsetI)
apply (erule CollectE conjE)+
apply (erule alg_F1set)

```

```

apply (erule subset_trans)
apply (rule UN_least)
apply (erule allE)
apply (drule mp)
  apply (erule underS_E)
apply (drule mp)
  apply (erule underS_Field)
apply (erule conjE)+
apply assumption

```

```

apply (erule subset_trans)
apply (rule UN_least)
apply (erule allE)
apply (drule mp)
  apply (erule underS_E)
apply (drule mp)
  apply (erule underS_Field)
apply (erule conjE)+
apply assumption

```

```

apply (rule ord_eq_le_trans)
apply (erule min_algs2)
apply (rule Un_least)
apply (rule UN_least)
apply (erule allE)
apply (drule mp)
  apply (erule underS_E)
apply (drule mp)
  apply (erule underS_Field)
apply (erule conjE)+
apply assumption
apply (rule image_subsetI)
apply (erule CollectE conjE)+
apply (erule alg_F2set)

```

```

apply (erule subset_trans)
apply (rule UN_least)
apply (erule allE)
apply (drule mp)
  apply (erule underS_E)
apply (drule mp)
  apply (erule underS_Field)
apply (erule conjE)+
apply assumption

```

```

apply (erule subset_trans)
apply (rule UN_least)
apply (erule allE)
apply (drule mp)
  apply (erule underS_E)
apply (drule mp)
  apply (erule underS_Field)
apply (erule conjE)+
apply assumption
done

```

```

lemma least_min_alg1: alg B1 B2 s1 s2  $\implies$  min_alg1 s1 s2  $\subseteq$  B1
apply (rule ord_eq_le_trans[OF min_alg1_def])
apply (rule UN_least)

```

```

apply (drule least_min_algs)
apply (drule mp)
  apply assumption
apply (erule conjE)+
apply assumption
done

```

```

lemma least_min_alg2: alg B1 B2 s1 s2  $\implies$  min_alg2 s1 s2  $\subseteq$  B2
apply (rule ord_eq_le_trans[OF min_alg2_def])
apply (rule UN_least)
apply (drule least_min_algs)
apply (drule mp)
  apply assumption
apply (erule conjE)+
apply assumption
done

```

```

lemma mor_incl_min_alg:
  alg B1 B2 s1 s2  $\implies$ 
  mor (min_alg1 s1 s2) (min_alg2 s1 s2) s1 s2 B1 B2 s1 s2 id id
apply (rule mor_incl)
  apply (erule least_min_alg1)
  apply (erule least_min_alg2)
done

```

1.5 Initiality

The following "happens" to be the type (for our particular construction) of the initial algebra carrier:

```

type-synonym 'a1 F1init_type = ('a1, 'a1 ASucFbd_type, 'a1 ASucFbd_type) F1  $\Rightarrow$  'a1 ASucFbd_type
type-synonym 'a1 F2init_type = ('a1, 'a1 ASucFbd_type, 'a1 ASucFbd_type) F2  $\Rightarrow$  'a1 ASucFbd_type

```

```

typedef 'a1 IIT =
  UNIV ::
  (('a1 ASucFbd_type set  $\times$  'a1 ASucFbd_type set)  $\times$  ('a1 F1init_type  $\times$  'a1 F2init_type)) set
by (rule exI) (rule UNIV_I)

```

1.6 Initial Algebras

abbreviation II :: 'a1 IIT set **where**

```

  II  $\equiv$  {Abs_IIT ((B1, B2), (s1, s2)) | B1 B2 s1 s2. alg B1 B2 s1 s2}

```

definition str_init1 **where**

```

  str_init1 (dummy :: 'a1)
  (y::('a1, 'a1 IIT  $\Rightarrow$  'a1 ASucFbd_type, 'a1 IIT  $\Rightarrow$  'a1 ASucFbd_type) F1)
  (i :: 'a1 IIT) =
  fst (snd (Rep_IIT i))
  (F1map id ( $\lambda$ f :: 'a1 IIT  $\Rightarrow$  'a1 ASucFbd_type. f i) ( $\lambda$ f. f i) y)

```

definition str_init2 **where**

```

  str_init2 (dummy :: 'a1) y (i :: 'a1 IIT) =
  snd (snd (Rep_IIT i)) (F2map id ( $\lambda$ f. f i) ( $\lambda$ f. f i) y)

```

abbreviation car_init1 **where**

```

  car_init1 dummy  $\equiv$  min_alg1 (str_init1 dummy) (str_init2 dummy)

```

abbreviation car_init2 **where**

```

  car_init2 dummy  $\equiv$  min_alg2 (str_init1 dummy) (str_init2 dummy)

```

lemma alg_select:

```

 $\forall i \in II.$  alg (fst (fst (Rep_IIT i))) (snd (fst (Rep_IIT i)))
  (fst (snd (Rep_IIT i))) (snd (snd (Rep_IIT i)))

```

```

apply (rule ballI)
apply (erule CollectE exE conjE)+
apply (tactic  $\langle\langle$  hyp_subst_tac @{context} 1  $\rangle\rangle$ )
unfolding fst_conv_snd_conv Abs_IIT_inverse[OF UNIV_I]
apply assumption
done

```

```

lemma mor_select:
  [[i ∈ II;
    mor (fst (fst (Rep_IIT i))) (snd (fst (Rep_IIT i)))
      (fst (snd (Rep_IIT i))) (snd (snd (Rep_IIT i))) UNIV UNIV s1' s2' f g]] ==>
  mor (car_init1 dummy) (car_init2 dummy) (str_init1 dummy) (str_init2 dummy) UNIV UNIV s1' s2' (f ∘ (λh.
h i)) (g ∘ (λh. h i))
apply (rule mor_cong)
  apply (rule sym)
  apply (rule o_id)
  apply (rule sym)
  apply (rule o_id)
apply (tactic << rtac @{context} (Thm.permute_premis 0 1 @{thm mor_comp}) 1 >>)
apply (tactic << etac @{context} (Thm.permute_premis 0 1 @{thm mor_comp}) 1 >>)
apply (tactic << rtac @{context} (@{thm mor_def} RS iffD2) 1 >>)
apply (rule conjI)

  apply (rule conjI)
  apply (rule ballI)
  apply (erule bspec[rotated])
  apply (erule CollectE)
  apply assumption

  apply (rule ballI)
  apply (erule bspec[rotated])
  apply (erule CollectE)
  apply assumption

  apply (rule conjI)
  apply (rule ballI)
  apply (rule str_init1_def)

  apply (rule ballI)
  apply (rule str_init2_def)

apply (rule mor_incl_min_alg)

apply (erule thin_rl)+
apply (tactic << rtac @{context} (@{thm alg_def} RS iffD2) 1 >>)
apply (rule conjI)
apply (rule ballI)
apply (erule CollectE conjE)+
apply (rule CollectI)
apply (rule ballI)
apply (erule bspec[OF alg_select])
apply (rule ssubst_mem[OF str_init1_def])
apply (erule alg_F1set)

  apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(2))
  apply (rule subset_trans)
  apply (erule image_mono)
  apply (rule image_Collect_subsetI)
  apply (erule bspec)
  apply assumption

  apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(3))
  apply (rule subset_trans)
  apply (erule image_mono)
  apply (rule image_Collect_subsetI)
  apply (erule bspec)
  apply assumption

```

```

apply (rule ballI)
apply (erule CollectE conjE)+
apply (rule CollectI)
apply (rule ballI)
apply (frule bspec[OF alg_select])
apply (rule ssubst_mem[OF str_init2_def])
apply (erule alg_F2set)

```

```

apply (rule ord_eq_le_trans)
apply (rule F2.set_map(2))
apply (rule subset_trans)
apply (erule image_mono)
apply (rule image_Collect_subsetI)
apply (erule bspec)
apply assumption

```

```

apply (rule ord_eq_le_trans)
apply (rule F2.set_map(3))
apply (rule subset_trans)
apply (erule image_mono)
apply (rule image_Collect_subsetI)
apply (erule bspec)
apply assumption
done

```

lemma *init_unique_mor*:

```

[[a1 ∈ car_init1 dummy; a2 ∈ car_init2 dummy;
  mor (car_init1 dummy) (car_init2 dummy) (str_init1 dummy) (str_init2 dummy) B1 B2 s1 s2 f1 f2;
  mor (car_init1 dummy) (car_init2 dummy) (str_init1 dummy) (str_init2 dummy) B1 B2 s1 s2 g1 g2]] ⇒
f1 a1 = g1 a1 ∧ f2 a2 = g2 a2

```

```

apply (rule conjI)
apply (erule prop_restrict)
apply (erule thin_rl)
apply (rule least_min_alg1)
apply (tactic ⟨⟨ rtac @{context} (@{thm alg_def} RS iffD2) 1 ⟩⟩)
apply (rule conjI)
apply (rule ballI)
apply (rule CollectI)
apply (erule CollectE conjE)+
apply (rule conjI)

```

```

apply (rule alg_F1set[OF alg_min_alg])
apply (erule subset_trans)
apply (rule Collect_restrict)
apply (erule subset_trans)
apply (rule Collect_restrict)

```

```

apply (rule trans)
apply (erule morE1)
apply (rule set_mp)
apply (rule Fin_mono23)
apply (rule Collect_restrict)
apply (rule Collect_restrict)
apply (rule CollectI)
apply (rule conjI)
apply assumption
apply (rule conjI)
apply assumption
apply assumption

```

```

apply (rule trans)

```

apply (rule arg_cong[OF F1.map_cong0])
apply (rule refl)
apply (erule prop_restrict)
apply assumption
apply (erule prop_restrict)
apply assumption

apply (rule sym)
apply (erule morE1)
apply (rule set_mp)
apply (rule F1in_mono23)
apply (rule Collect_restrict)
apply (rule Collect_restrict)
apply (rule CollectI)
apply (rule conjI)
apply assumption
apply (rule conjI)
apply assumption
apply assumption

apply (rule ballI)
apply (rule CollectI)
apply (erule CollectE conjE)+
apply (rule conjI)

apply (rule alg_F2set[OF alg_min_alg])
apply (erule subset_trans)
apply (rule Collect_restrict)
apply (erule subset_trans)
apply (rule Collect_restrict)

apply (rule trans)
apply (erule morE2)
apply (rule set_mp)
apply (rule F2in_mono23)
apply (rule Collect_restrict)
apply (rule Collect_restrict)
apply (rule CollectI)
apply (rule conjI)
apply assumption
apply (rule conjI)
apply assumption
apply assumption

apply (rule trans)
apply (rule arg_cong[OF F2.map_cong0])
apply (rule refl)
apply (erule prop_restrict)
apply assumption
apply (erule prop_restrict)
apply assumption

apply (rule sym)
apply (erule morE2)
apply (rule set_mp)
apply (rule F2in_mono23)
apply (rule Collect_restrict)
apply (rule Collect_restrict)
apply (rule CollectI)
apply (rule conjI)
apply assumption
apply (rule conjI)
apply assumption

apply *assumption*

apply (*erule thin_rl*)
apply (*erule prop_restrict*)
apply (*rule least_min_alg2*)
apply (*tactic* \ll *rtac* $\@_{\{context\}}$ ($\@_{\{thm\ alg_def\}}$ *RS iffD2*) 1 \gg)
apply (*rule conjI*)
apply (*rule ballI*)
apply (*rule CollectI*)
apply (*erule CollectE conjE*)
apply (*rule conjI*)

apply (*rule alg_F1set*[*OF alg_min_alg*])
apply (*erule subset_trans*)
apply (*rule Collect_restrict*)
apply (*erule subset_trans*)
apply (*rule Collect_restrict*)

apply (*rule trans*)
apply (*erule morE1*)
apply (*rule set_mp*)
apply (*rule F1in_mono23*)
apply (*rule Collect_restrict*)
apply (*rule Collect_restrict*)
apply (*rule CollectI*)
apply (*rule conjI*)
apply *assumption*
apply (*rule conjI*)
apply *assumption*
apply *assumption*

apply (*rule trans*)
apply (*rule arg_cong*[*OF F1.map_cong0*])
apply (*rule refl*)
apply (*erule prop_restrict*)
apply *assumption*
apply (*erule prop_restrict*)
apply *assumption*

apply (*rule sym*)
apply (*erule morE1*)
apply (*rule set_mp*)
apply (*rule F1in_mono23*)
apply (*rule Collect_restrict*)
apply (*rule Collect_restrict*)
apply (*rule CollectI*)
apply (*rule conjI*)
apply *assumption*
apply (*rule conjI*)
apply *assumption*
apply *assumption*

apply (*rule ballI*)
apply (*rule CollectI*)
apply (*erule CollectE conjE*)
apply (*rule conjI*)

apply (*rule alg_F2set*[*OF alg_min_alg*])
apply (*erule subset_trans*)
apply (*rule Collect_restrict*)
apply (*erule subset_trans*)
apply (*rule Collect_restrict*)

```

apply (rule trans)
apply (erule morE2)
apply (rule set_mp)
  apply (rule F2in_mono23)
    apply (rule Collect_restrict)
  apply (rule Collect_restrict)
apply (rule CollectI)
apply (rule conjI)
apply assumption
apply (rule conjI)
apply assumption
apply assumption

```

```

apply (rule trans)
apply (rule arg_cong[OF F2.map_cong0])
  apply (rule refl)
apply (erule prop_restrict)
apply assumption
apply (erule prop_restrict)
apply assumption

```

```

apply (rule sym)
apply (erule morE2)
apply (rule set_mp)
  apply (rule F2in_mono23)
    apply (rule Collect_restrict)
  apply (rule Collect_restrict)
apply (rule CollectI)
apply (rule conjI)
apply assumption
apply (rule conjI)
apply assumption
apply assumption
done

```

abbreviation closed where

$$\begin{aligned}
 \text{closed dummy } \phi_1 \phi_2 \equiv & ((\forall x \in F1in \text{ UNIV } (\text{car_init1 dummy}) (\text{car_init2 dummy}). \\
 & (\forall z \in F1set2 \ x. \phi_1 z) \wedge (\forall z \in F1set3 \ x. \phi_2 z) \longrightarrow \phi_1 (\text{str_init1 dummy } x)) \wedge \\
 & (\forall x \in F2in \text{ UNIV } (\text{car_init1 dummy}) (\text{car_init2 dummy}). \\
 & (\forall z \in F2set2 \ x. \phi_1 z) \wedge (\forall z \in F2set3 \ x. \phi_2 z) \longrightarrow \phi_2 (\text{str_init2 dummy } x)))
 \end{aligned}$$

lemma *init_induct*: *closed dummy* $\phi_1 \phi_2 \implies$

```

 $(\forall x \in \text{car\_init1 dummy. } \phi_1 x) \wedge (\forall x \in \text{car\_init2 dummy. } \phi_2 x)$ 
apply (rule conjI)
apply (rule ballI)
apply (erule prop_restrict)
apply (rule least_min_alg1)
apply (tactic << rtac @{context} (@{thm alg_def} RS iffD2) 1 >>)

```

```

apply (rule conjI)
apply (rule ballI)
apply (rule CollectI)
apply (erule CollectE conjE)+
apply (rule conjI)

```

```

apply (rule alg_F1set[OF alg_min_alg])
apply (erule subset_trans)
apply (rule Collect_restrict)
apply (erule subset_trans)
apply (rule Collect_restrict)

```

```

apply (rule mp)

```



```

apply (erule bspec)
apply (rule CollectI)
apply (rule conjI)
  apply assumption
apply (rule conjI)
  apply (erule subset_trans)
  apply (rule Collect_restrict)
apply (erule subset_trans)
apply (rule Collect_restrict)

apply (rule conjI)
apply (rule ballI)
apply (erule prop_restrict)
apply assumption
apply (rule ballI)
apply (erule prop_restrict)
apply assumption

apply (rule ballI)
apply (rule CollectI)
apply (erule CollectE conjE)+
apply (rule conjI)

apply (rule alg_F2set[OF alg_min_alg])
  apply (erule subset_trans)
  apply (rule Collect_restrict)
apply (erule subset_trans)
apply (rule Collect_restrict)

apply (rule mp)
apply (erule bspec)
apply (rule CollectI)
apply (rule conjI)
  apply assumption
apply (rule conjI)
  apply (erule subset_trans)
  apply (rule Collect_restrict)
apply (erule subset_trans)
apply (rule Collect_restrict)

apply (rule conjI)
apply (rule ballI)
apply (erule prop_restrict)
apply assumption
apply (rule ballI)
apply (erule prop_restrict)
apply assumption

apply (rule ballI)
apply (erule prop_restrict)
apply (rule least_min_alg2)
apply (tactic << rtac @{context} (@{thm alg_def} RS iffD2) 1 >>)

apply (rule conjI)
apply (rule ballI)
apply (rule CollectI)
apply (erule CollectE conjE)+
apply (rule conjI)

apply (rule alg_F1set[OF alg_min_alg])
  apply (erule subset_trans)
  apply (rule Collect_restrict)

```

```

apply (erule subset_trans)
apply (rule Collect_restrict)

```

```

apply (rule mp)
apply (erule bspec)
apply (rule CollectI)
apply (rule conjI)
apply assumption
apply (rule conjI)
apply (erule subset_trans)
apply (rule Collect_restrict)
apply (erule subset_trans)
apply (rule Collect_restrict)

```

```

apply (rule conjI)
apply (rule ballI)
apply (erule prop_restrict)
apply assumption
apply (rule ballI)
apply (erule prop_restrict)
apply assumption

```

```

apply (rule ballI)
apply (rule CollectI)
apply (erule CollectE conjE)+
apply (rule conjI)

```

```

apply (rule alg_F2set[OF alg_min_alg])
apply (erule subset_trans)
apply (rule Collect_restrict)
apply (erule subset_trans)
apply (rule Collect_restrict)

```

```

apply (rule mp)
apply (erule bspec)
apply (rule CollectI)
apply (rule conjI)
apply assumption
apply (rule conjI)
apply (erule subset_trans)
apply (rule Collect_restrict)
apply (erule subset_trans)
apply (rule Collect_restrict)

```

```

apply (rule conjI)
apply (rule ballI)
apply (erule prop_restrict)
apply assumption
apply (rule ballI)
apply (erule prop_restrict)
apply assumption
done

```

1.7 The datatype

```

typedef (overloaded) 'a1 IF1 = car_init1 (undefined :: 'a1)
apply (rule iffD2)
apply (rule ex_in_conv)
apply (rule conjunct1)
apply (rule alg_not_empty)
apply (rule alg_min_alg)
done

```

```

typedef (overloaded) 'a1 IF2 = car_init2 (undefined :: 'a1)
  apply (rule iffD2)
  apply (rule ex_in_conv)
  apply (rule conjunct2)
  apply (rule alg_not_empty)
  apply (rule alg_min_alg)
done

```

```

definition ctor1 where ctor1 = Abs_IF1 o str_init1 undefined o F1map id Rep_IF1 Rep_IF2
definition ctor2 where ctor2 = Abs_IF2 o str_init2 undefined o F2map id Rep_IF1 Rep_IF2

```

lemma mor_Rep_IF:

```

mor (UNIV :: 'a IF1 set) (UNIV :: 'a IF2 set) ctor1 ctor2
  (car_init1 undefined) (car_init2 undefined) (str_init1 undefined) (str_init2 undefined) Rep_IF1 Rep_IF2

```

unfolding mor_def ctor1_def ctor2_def o_ apply

```

apply (rule conjI)
apply (rule conjI)
  apply (rule ballI)
  apply (rule Rep_IF1)
apply (rule ballI)
apply (rule Rep_IF2)

```

```

apply (rule conjI)
apply (rule ballI)
apply (rule Abs_IF1_inverse)
apply (rule alg_F1set[OF alg_min_alg])
  apply (rule ord_eq_le_trans[OF F1.set_map(2)])
  apply (rule image_subsetI)
  apply (rule Rep_IF1)
apply (rule ord_eq_le_trans[OF F1.set_map(3)])
apply (rule image_subsetI)
apply (rule Rep_IF2)

```

```

apply (rule ballI)
apply (rule Abs_IF2_inverse)
apply (rule alg_F2set[OF alg_min_alg])
  apply (rule ord_eq_le_trans[OF F2.set_map(2)])
  apply (rule image_subsetI)
  apply (rule Rep_IF1)
apply (rule ord_eq_le_trans[OF F2.set_map(3)])
apply (rule image_subsetI)
apply (rule Rep_IF2)
done

```

lemma mor_Abs_IF:

```

mor (car_init1 undefined) (car_init2 undefined)
  (str_init1 undefined) (str_init2 undefined) UNIV UNIV ctor1 ctor2 Abs_IF1 Abs_IF2

```

unfolding mor_def ctor1_def ctor2_def o_ apply

```

apply (rule conjI)
apply (rule conjI)
  apply (rule ballI)
  apply (rule UNIV_I)
apply (rule ballI)
apply (rule UNIV_I)

```

```

apply (rule conjI)
apply (rule ballI)
apply (erule CollectE conjE)+
apply (rule sym[OF arg_cong[OF trans[OF F1map_comp_id F1map_congL]]])
  apply (rule ballI[OF trans[OF o_apply]])
  apply (erule Abs_IF1_inverse[OF set_mp])
  apply assumption
apply (rule ballI[OF trans[OF o_apply]])

```

```

apply (erule Abs_IF2_inverse[OF set_mp])
apply assumption

```

```

apply (rule ballI)
apply (erule CollectE conjE)+
apply (rule sym[OF arg_cong[OF trans[OF F2map_comp_id F2map_congL]]])
  apply (rule ballI[OF trans[OF o_apply]])
  apply (erule Abs_IF1_inverse[OF set_mp])
  apply assumption
apply (rule ballI[OF trans[OF o_apply]])
apply (erule Abs_IF2_inverse[OF set_mp])
apply assumption
done

```

lemma copy:

```

[[alg B1 B2 s1 s2; bij_betw f B1' B1; bij_betw g B2' B2]] ==>
  ∃ f' g'. alg B1' B2' f' g' ∧ mor B1' B2' f' g' B1 B2 s1 s2 f g
apply (rule exI)+
apply (rule conjI)
apply (tactic ⟨⟨ rtac @ {context} (@ {thm alg_def} RS iffD2) 1 ⟩⟩)
apply (rule conjI)
apply (rule ballI)
apply (erule CollectE conjE)+
apply (rule set_mp)
  apply (rule equalityD1)
  apply (erule bij_betw_imp_surj_on[OF bij_betw_the_inv_into])
apply (rule imageI)
apply (erule alg_F1set)
  apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(2))
  apply (rule subset_trans)
  apply (erule image_mono)
  apply (rule equalityD1)
  apply (erule bij_betw_imp_surj_on)
apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(3))
apply (rule subset_trans)
  apply (erule image_mono)
apply (rule equalityD1)
apply (erule bij_betw_imp_surj_on)

apply (rule ballI)
apply (erule CollectE conjE)+
apply (rule set_mp)
  apply (rule equalityD1)
  apply (erule bij_betw_imp_surj_on[OF bij_betw_the_inv_into])
apply (rule imageI)
apply (erule alg_F2set)
  apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(2))
  apply (rule subset_trans)
  apply (erule image_mono)
  apply (rule equalityD1)
  apply (erule bij_betw_imp_surj_on)
apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(3))
apply (rule subset_trans)
  apply (erule image_mono)
apply (rule equalityD1)
apply (erule bij_betw_imp_surj_on)

apply (tactic ⟨⟨ rtac @ {context} (@ {thm mor_def} RS iffD2) 1 ⟩⟩)
apply (rule conjI)

```

```

apply (rule conjI)
apply (erule bij_betwE)
apply (erule bij_betwE)

```

```

apply (rule conjI)
apply (rule ballI)
apply (erule CollectE conjE)+
apply (erule f_the_inv_into_f_bij_betw)
apply (erule alg_F1set)
apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(2))
apply (rule subset_trans)
  apply (erule image_mono)
apply (rule equalityD1)
apply (erule bij_betw_imp_surj_on)
apply (rule ord_eq_le_trans)
apply (rule F1.set_map(3))
apply (rule subset_trans)
apply (erule image_mono)
apply (rule equalityD1)
apply (erule bij_betw_imp_surj_on)

```

```

apply (rule ballI)
apply (erule CollectE conjE)+
apply (erule f_the_inv_into_f_bij_betw)
apply (erule alg_F2set)
apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(2))
apply (rule subset_trans)
  apply (erule image_mono)
apply (rule equalityD1)
apply (erule bij_betw_imp_surj_on)
apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(3))
apply (rule subset_trans)
  apply (erule image_mono)
apply (rule equalityD1)
apply (erule bij_betw_imp_surj_on)
done

```

lemma *init_ex_mor*:

```

 $\exists f g. \text{mor } UNIV \ UNIV \ \text{ctor1 } \text{ctor2 } \ UNIV \ UNIV \ s1 \ s2 \ f \ g$ 
apply (insert ex_bij_betw[OF card_of_min_alg1, of s1 s2]
  ex_bij_betw[OF card_of_min_alg2, of s1 s2])
apply (erule exE)+
apply (rule rev_mp)
apply (rule copy[OF alg_min_alg])
  apply assumption
  apply assumption
apply (rule impI)
apply (erule exE conjE)+

apply (rule exI)+
apply (rule mor_comp)
  apply (rule mor_Rep_IF)
apply (rule mor_select)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply assumption
unfolding fst_conv snd_conv Abs_IIT_inverse[OF UNIV_I]
apply (erule mor_comp)

```

```

apply (rule mor_incl)
apply (rule subset_UNIV)
apply (rule subset_UNIV)
done

```

Iteration

abbreviation *fold* **where**

```

fold s1 s2  $\equiv$  (SOME f. mor UNIV UNIV ctor1 ctor2 UNIV UNIV s1 s2 (fst f) (snd f))

```

definition *fold1* **where** *fold1 s1 s2 = fst (fold s1 s2)*

definition *fold2* **where** *fold2 s1 s2 = snd (fold s1 s2)*

lemma *mor_fold*:

```

mor UNIV UNIV ctor1 ctor2 UNIV UNIV s1 s2 (fold1 s1 s2) (fold2 s1 s2)

```

```

unfolding fold1_def fold2_def

```

```

apply (rule rev_mp)

```

```

apply (rule init_ex_mor)

```

```

apply (rule impI)

```

```

apply (erule exE)

```

```

apply (erule exE)

```

```

apply (rule someI[of %(f :: ('a IF1  $\Rightarrow$  'b)  $\times$  ('a IF2  $\Rightarrow$  'c)).

```

```

mor UNIV UNIV ctor1 ctor2 UNIV UNIV s1 s2 (fst f) (snd f]])

```

```

apply (erule mor_cong[OF fst_conv snd_conv])

```

```

done

```

ML $\langle\langle$

```

val fold1 = rule_by_tactic @{context}

```

```

  (rtac @{context} CollectI 1 THEN BNF_Util.CONJ_WRAP (K (rtac @{context} @{thm subset_UNIV} 1))

```

```

(1 upto 3))

```

```

  @{thm morE1[OF mor_fold]}

```

```

val fold2 = rule_by_tactic @{context}

```

```

  (rtac @{context} CollectI 1 THEN BNF_Util.CONJ_WRAP (K (rtac @{context} @{thm subset_UNIV} 1))

```

```

(1 upto 3))

```

```

  @{thm morE2[OF mor_fold]}

```

$\rangle\rangle$

theorem *fold1*:

```

(fold1 s1 s2) (ctor1 x) = s1 (F1map id (fold1 s1 s2) (fold2 s1 s2) x)

```

```

apply (rule morE1)

```

```

apply (rule mor_fold)

```

```

apply (rule CollectI)

```

```

apply (rule conjI)

```

```

apply (rule subset_UNIV)

```

```

apply (rule conjI)

```

```

apply (rule subset_UNIV)

```

```

apply (rule subset_UNIV)

```

```

done

```

theorem *fold2*:

```

(fold2 s1 s2) (ctor2 x) = s2 (F2map id (fold1 s1 s2) (fold2 s1 s2) x)

```

```

apply (rule morE2)

```

```

apply (rule mor_fold)

```

```

apply (rule CollectI)

```

```

apply (rule conjI)

```

```

apply (rule subset_UNIV)

```

```

apply (rule conjI)

```

```

apply (rule subset_UNIV)

```

```

apply (rule subset_UNIV)

```

```

done

```

lemma *mor_UNIV*: mor UNIV UNIV *s1 s2* UNIV UNIV *s1' s2'* *f g* \longleftrightarrow

```

f o s1 = s1' o F1map id f g  $\wedge$  g o s2 = s2' o F2map id f g

```

```

apply (rule iffI)
apply (rule conjI)
apply (rule ext)
apply (rule trans)
  apply (rule o_apply)
apply (rule trans)
apply (erule morE1)
apply (rule CollectI)
apply (rule conjI)
  apply (rule subset_UNIV)
apply (rule conjI)
  apply (rule subset_UNIV)
apply (rule subset_UNIV)
apply (rule sym[OF o_apply])

```

```

apply (rule ext)
apply (rule trans)
apply (rule o_apply)
apply (rule trans)
apply (erule morE2)
apply (rule CollectI)
apply (rule conjI)
  apply (rule subset_UNIV)
apply (rule conjI)
  apply (rule subset_UNIV)
apply (rule subset_UNIV)
apply (rule sym[OF o_apply])

```

```

apply (tactic << rtac @{context} (@{thm mor_def} RS iffD2) 1 >>)
apply (rule conjI)
apply (rule conjI)
  apply (rule ballI)
  apply (rule UNIV_I)
apply (rule ballI)
apply (rule UNIV_I)
apply (erule conjE)
apply (erule iffD1[OF fun_eq_iff])
apply (erule iffD1[OF fun_eq_iff])
apply (rule conjI)
apply (rule ballI)
apply (erule allE)+
apply (rule trans)
  apply (erule trans[OF sym[OF o_apply]])
apply (rule o_apply)
apply (rule ballI)
apply (erule allE)+
apply (rule trans)
  apply (erule trans[OF sym[OF o_apply]])
apply (rule o_apply)
done

```

lemma fold_unique_mor: mor UNIV UNIV ctor1 ctor2 UNIV UNIV s1 s2 f g \implies
 $f = \text{fold1 } s1 \ s2 \wedge g = \text{fold2 } s1 \ s2$

```

apply (rule conjI)
apply (rule surj_fun_eq)
  apply (rule type_definition.Abs_image[OF type_definition_IF1])
apply (rule ballI)
apply (rule conjunct1)
apply (rule init_unique_mor)
  apply assumption
  apply (rule Rep_IF2)
apply (rule mor_comp)
apply (rule mor_Abs_IF)

```

```

apply assumption
apply (rule mor_comp)
apply (rule mor_Abs_IF)
apply (rule mor_fold)

apply (rule surj_fun_eq)
apply (rule type_definition.Abs_image[OF type_definition_IF2])
apply (rule ballI)
apply (rule conjunct2)
apply (rule init_unique_mor)
  apply (rule Rep_IF1)
  apply assumption
apply (rule mor_comp)
apply (rule mor_Abs_IF)
apply assumption
apply (rule mor_comp)
apply (rule mor_Abs_IF)
apply (rule mor_fold)
done

lemmas fold_unique = fold_unique_mor[OF iffD2[OF mor_UNIV], OF conjI]

lemmas fold1_ctor = sym[OF conjunct1[OF fold_unique_mor[OF mor_incl[OF subset_UNIV subset_UNIV]]]]
lemmas fold2_ctor = sym[OF conjunct2[OF fold_unique_mor[OF mor_incl[OF subset_UNIV subset_UNIV]]]]

Case distinction

lemmas ctor1_o_fold1 =
  trans[OF conjunct1[OF fold_unique_mor[OF mor_comp[OF mor_fold mor_str]]] fold1_ctor]
lemmas ctor2_o_fold2 =
  trans[OF conjunct2[OF fold_unique_mor[OF mor_comp[OF mor_fold mor_str]]] fold2_ctor]

definition dtor1 = fold1 (F1map id ctor1 ctor2) (F2map id ctor1 ctor2)
definition dtor2 = fold2 (F1map id ctor1 ctor2) (F2map id ctor1 ctor2)

ML  $\ll$  Local_Defs.fold @ {context} @ {thms dtor1_def} @ {thm ctor1_o_fold1}  $\gg$ 
ML  $\ll$  Local_Defs.fold @ {context} @ {thms dtor2_def} @ {thm ctor2_o_fold2}  $\gg$ 

lemma ctor1_o_dtor1: ctor1 o dtor1 = id
  unfolding dtor1_def
  apply (rule ctor1_o_fold1)
done

lemma ctor2_o_dtor2: ctor2 o dtor2 = id
  unfolding dtor2_def
  apply (rule ctor2_o_fold2)
done

lemma dtor1_o_ctor1: dtor1 o ctor1 = id
  apply (rule ext)
  apply (rule trans[OF o_apply])
  apply (rule trans[OF fun_cong[OF dtor1_def]])
  apply (rule trans[OF fold1])
  apply (rule trans[OF F1map_comp_id])
  apply (rule trans[OF F1map_congL])
  apply (rule ballI)
  apply (rule trans[OF fun_cong[OF ctor1_o_fold1] id_apply])
  apply (rule ballI)
  apply (rule trans[OF fun_cong[OF ctor2_o_fold2] id_apply])
  apply (rule sym[OF id_apply])
done

lemma dtor2_o_ctor2: dtor2 o ctor2 = id
  apply (rule ext)

```



```

apply (rule trans[OF o_apply])
apply (rule trans[OF fun_cong[OF dtor2_def]])
apply (rule trans[OF fold2])
apply (rule trans[OF F2map_comp_id])
apply (rule trans[OF F2map_congL])
  apply (rule ballI)
  apply (rule trans[OF fun_cong[OF ctor1_o_fold1] id_apply])
  apply (rule ballI)
  apply (rule trans[OF fun_cong[OF ctor2_o_fold2] id_apply])
apply (rule sym[OF id_apply])
done

```

```

lemmas dtor1_ctor1 = pointfree_idE[OF dtor1_o_ctor1]
lemmas dtor2_ctor2 = pointfree_idE[OF dtor2_o_ctor2]
lemmas ctor1_dtor1 = pointfree_idE[OF ctor1_o_dtor1]
lemmas ctor2_dtor2 = pointfree_idE[OF ctor2_o_dtor2]

```

```

lemmas bij_dtor1 = o_bij[OF ctor1_o_dtor1 dtor1_o_ctor1]
lemmas inj_dtor1 = bij_is_inj[OF bij_dtor1]
lemmas surj_dtor1 = bij_is_surj[OF bij_dtor1]
lemmas dtor1_nchotomy = surjD[OF surj_dtor1]
lemmas dtor1_diff = inj_eq[OF inj_dtor1]
lemmas dtor1_cases = exE[OF dtor1_nchotomy]
lemmas bij_dtor2 = o_bij[OF ctor2_o_dtor2 dtor2_o_ctor2]
lemmas inj_dtor2 = bij_is_inj[OF bij_dtor2]
lemmas surj_dtor2 = bij_is_surj[OF bij_dtor2]
lemmas dtor2_nchotomy = surjD[OF surj_dtor2]
lemmas dtor2_diff = inj_eq[OF inj_dtor2]
lemmas dtor2_cases = exE[OF dtor2_nchotomy]

```

```

lemmas bij_ctor1 = o_bij[OF dtor1_o_ctor1 ctor1_o_dtor1]
lemmas inj_ctor1 = bij_is_inj[OF bij_ctor1]
lemmas surj_ctor1 = bij_is_surj[OF bij_ctor1]
lemmas ctor1_nchotomy = surjD[OF surj_ctor1]
lemmas ctor1_diff = inj_eq[OF inj_ctor1]
lemmas ctor1_cases = exE[OF ctor1_nchotomy]
lemmas bij_ctor2 = o_bij[OF dtor2_o_ctor2 ctor2_o_dtor2]
lemmas inj_ctor2 = bij_is_inj[OF bij_ctor2]
lemmas surj_ctor2 = bij_is_surj[OF bij_ctor2]
lemmas ctor2_nchotomy = surjD[OF surj_ctor2]
lemmas ctor2_diff = inj_eq[OF inj_ctor2]
lemmas ctor2_cases = exE[OF ctor2_nchotomy]

```

Primitive recursion

definition rec1 **where**

```

  rec1 s1 s2 = snd o fold1 (<ctor1 o F1map id fst fst, s1>) (<ctor2 o F2map id fst fst, s2>)

```

definition rec2 **where**

```

  rec2 s1 s2 = snd o fold2 (<ctor1 o F1map id fst fst, s1>) (<ctor2 o F2map id fst fst, s2>)

```

lemma fold1_o_ctor1: fold1 s1 s2 o ctor1 = s1 o F1map id (fold1 s1 s2) (fold2 s1 s2)

by (tactic ⟨rtac @{{context}} (BNF_Tactics.mk_pointfree2 @{{context}} @{{thm fold1}}) 1⟩)

lemma fold2_o_ctor2: fold2 s1 s2 o ctor2 = s2 o F2map id (fold1 s1 s2) (fold2 s1 s2)

by (tactic ⟨rtac @{{context}} (BNF_Tactics.mk_pointfree2 @{{context}} @{{thm fold2}}) 1⟩)

lemmas fst_rec1_pair =

```

  trans[OF conjunct1[OF fold_unique[OF
    trans[OF o_assoc[symmetric] trans[OF arg_cong2[of _ _ _ _ (o), OF refl
      trans[OF fold1_o_ctor1 convol_o]]], OF trans[OF fst_convol]]
    trans[OF o_assoc[symmetric] trans[OF arg_cong2[of _ _ _ _ (o), OF refl
      trans[OF fold2_o_ctor2 convol_o]]], OF trans[OF fst_convol]]]]]
  fold1_ctor, unfolded F1.map_comp0[of id, unfolded id_o] F2.map_comp0[of id, unfolded id_o] o_assoc,
  OF refl refl]

```

lemmas fst_rec2_pair =

```

trans[OF conjunct2[OF fold_unique[OF
  trans[OF o_assoc[symmetric] trans[OF arg_cong2[of _ _ _ _ (o), OF refl
    trans[OF fold1_o_ctor1 convol_o]]], OF trans[OF fst_convol]]
  trans[OF o_assoc[symmetric] trans[OF arg_cong2[of _ _ _ _ (o), OF refl
    trans[OF fold2_o_ctor2 convol_o]]], OF trans[OF fst_convol]]]]]
fold2_ctor, unfolded F1.map_comp0[of id, unfolded id_o] F2.map_comp0[of id, unfolded id_o] o_assoc,
OF refl refl]

```

theorem *rec1*: $rec1\ s1\ s2\ (ctor1\ x) = s1\ (F1map\ id\ (<id,\ rec1\ s1\ s2>)\ (<id,\ rec2\ s1\ s2>)\ x)$
unfolding *rec1_def rec2_def o_apply fold1_snd_convol'*
convol_expand_snd[OF fst_rec1_pair] convol_expand_snd[OF fst_rec2_pair] ..

theorem *rec2*: $rec2\ s1\ s2\ (ctor2\ x) = s2\ (F2map\ id\ (<id,\ rec1\ s1\ s2>)\ (<id,\ rec2\ s1\ s2>)\ x)$
unfolding *rec1_def rec2_def o_apply fold2_snd_convol'*
convol_expand_snd[OF fst_rec1_pair] convol_expand_snd[OF fst_rec2_pair] ..

lemma *rec_unique*:

```

f o ctor1 = s1 o F1map id <id , f> <id , g> ==>
g o ctor2 = s2 o F2map id <id , f> <id , g> ==> f = rec1 s1 s2 ^ g = rec2 s1 s2
unfolding rec1_def rec2_def convol_expand_snd'[OF fst_rec1_pair] convol_expand_snd'[OF fst_rec2_pair]
apply (rule fold_unique)
apply (unfold convol_o_id_o o_id F1.map_comp0[symmetric] F2.map_comp0[symmetric]
F1.map_id0 F2.map_id0 o_assoc[symmetric] fst_convol)
apply (erule arg_cong2[of _ _ _ _ BNF_Def.convol, OF refl])
apply (erule arg_cong2[of _ _ _ _ BNF_Def.convol, OF refl])
done

```

Induction

theorem *ctor_induct*:

```

[[ $\bigwedge x. (\bigwedge a. a \in F1set2\ x \implies phi1\ a) \implies (\bigwedge a. a \in F1set3\ x \implies phi2\ a) \implies phi1\ (ctor1\ x);$ 
 $\bigwedge x. (\bigwedge a. a \in F2set2\ x \implies phi1\ a) \implies (\bigwedge a. a \in F2set3\ x \implies phi2\ a) \implies phi2\ (ctor2\ x)]] \implies$ 
 $phi1\ a \wedge phi2\ b$ 
apply (rule mp)

```

```

apply (rule impI)
apply (erule conjE)
apply (rule conjI)
apply (rule iffD1[OF arg_cong[OF Rep_IF1_inverse]])
apply (erule bspec[OF Rep_IF1])
apply (rule iffD1[OF arg_cong[OF Rep_IF2_inverse]])
apply (erule bspec[OF Rep_IF2])
apply (rule init_induct)

```

apply (rule *conjI*)

```

apply (drule asm_rl)
apply (erule thin_rl)
apply (rule ballI)
apply (rule impI)
apply (rule iffD2[OF arg_cong[OF morE1[OF mor_Abs_IF]])
apply assumption
apply (erule CollectE conjE)+
apply (drule meta_spec)
apply (drule meta_mp)
apply (rule iffD1[OF arg_cong[OF Rep_IF1_inverse]])
apply (erule bspec)
apply (drule set_rev_mp)
apply (rule equalityD1)
apply (rule F1.set_map(2))
apply (erule imageE)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule ssubst_mem[OF Abs_IF1_inverse])
apply (erule set_mp)

```

apply assumption
 apply assumption

apply (drule meta_mp)
 apply (rule iffD1[OF arg_cong[OF Rep_IF2_inverse]])
 apply (erule bspec)
 apply (drule set_rev_mp)
 apply (rule equalityD1)
 apply (rule F1.set_map(3))
 apply (erule imageE)
 apply (tactic << hyp_subst_tac @ {context} 1 >>)
 apply (rule ssubst_mem[OF Abs_IF2_inverse])
 apply (erule set_mp)
 apply assumption
 apply assumption

apply assumption

apply (erule thin_rl)
 apply (drule asm_rl)
 apply (rule ballI)
 apply (rule impI)
 apply (rule iffD2[OF arg_cong[OF morE2[OF mor_Abs_IF]]])
 apply assumption
 apply (erule CollectE conjE)+
 apply (drule meta_spec)
 apply (drule meta_mp)
 apply (rule iffD1[OF arg_cong[OF Rep_IF1_inverse]])
 apply (erule bspec)
 apply (drule set_rev_mp)
 apply (rule equalityD1)
 apply (rule F2.set_map(2))
 apply (erule imageE)
 apply (tactic << hyp_subst_tac @ {context} 1 >>)
 apply (rule ssubst_mem[OF Abs_IF1_inverse])
 apply (erule set_mp)
 apply assumption
 apply assumption

apply (drule meta_mp)
 apply (rule iffD1[OF arg_cong[OF Rep_IF2_inverse]])
 apply (erule bspec)
 apply (drule set_rev_mp)
 apply (rule equalityD1)
 apply (rule F2.set_map(3))
 apply (erule imageE)
 apply (tactic << hyp_subst_tac @ {context} 1 >>)
 apply (rule ssubst_mem[OF Abs_IF2_inverse])
 apply (erule set_mp)
 apply assumption
 apply assumption

apply assumption
 done

theorem ctor_induct2:

$$\llbracket \begin{aligned} & \forall x y. (\bigwedge a b. a \in F1set2\ x \implies b \in F1set2\ y \implies phi1\ a\ b) \implies \\ & (\bigwedge a b. a \in F1set3\ x \implies b \in F1set3\ y \implies phi2\ a\ b) \implies phi1\ (ctor1\ x)\ (ctor1\ y); \\ & \forall x y. (\bigwedge a b. a \in F2set2\ x \implies b \in F2set2\ y \implies phi1\ a\ b) \implies \\ & (\bigwedge a b. a \in F2set3\ x \implies b \in F2set3\ y \implies phi2\ a\ b) \implies phi2\ (ctor2\ x)\ (ctor2\ y) \rrbracket \implies \\ & phi1\ a1\ b1 \wedge phi2\ a2\ b2 \end{aligned}$$
 apply (rule rev_mp)
 apply (rule ctor_induct[of %a1. (∀ x. phi1 a1 x) %a2. (∀ y. phi2 a2 y) a1 a2])

```

apply (rule allI[OF conjunct1[OF ctor_induct[OF asm_rl TrueI]])
apply (drule meta_spec2)
apply (erule thin_rl)
apply (tactic << (dtac @ {context} @ {thm meta_mp} THEN_ALL_NEW Goal.norm_hhf_tac @ {context}) 1 >>)
  apply (drule meta_spec)+
  apply (erule meta_mp[OF spec])
  apply assumption
apply (drule meta_mp)
  apply (drule meta_spec)+
  apply (erule meta_mp[OF spec])
  apply assumption
apply assumption

apply (rule allI[OF conjunct2[OF ctor_induct[OF TrueI asm_rl]])
apply (erule thin_rl)
apply (drule meta_spec2)
apply (drule meta_mp)
  apply (drule meta_spec)+
  apply (erule meta_mp[OF spec])
  apply assumption
apply (erule meta_mp)
  apply (drule meta_spec)+
  apply (erule meta_mp[OF spec])
  apply assumption

apply (rule impI)
apply (erule conjE allE)+
apply (rule conjI)
  apply assumption
apply assumption
done

```

1.8 The Result as an BNF

The map operator

abbreviation *IF1map* **where** $IF1map\ f \equiv fold1\ (ctor1\ o\ (F1map\ f\ id\ id))\ (ctor2\ o\ (F2map\ f\ id\ id))$
abbreviation *IF2map* **where** $IF2map\ f \equiv fold2\ (ctor1\ o\ (F1map\ f\ id\ id))\ (ctor2\ o\ (F2map\ f\ id\ id))$

theorem *IF1map*:

```

(IF1map f) o ctor1 = ctor1 o (F1map f (IF1map f) (IF2map f))
apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule trans[OF fold1])
apply (rule trans[OF o_apply])
apply (rule trans[OF arg_cong[OF F1map_comp_id]])
apply (rule trans[OF arg_cong[OF F1.map_cong0]])
  apply (rule refl)
  apply (rule trans[OF o_apply])
  apply (rule id_apply)
apply (rule trans[OF o_apply])
apply (rule id_apply)
apply (rule sym[OF o_apply])
done

```

theorem *IF2map*:

```

(IF2map f) o ctor2 = ctor2 o (F2map f (IF1map f) (IF2map f))
apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule trans[OF fold2])
apply (rule trans[OF o_apply])
apply (rule trans[OF arg_cong[OF F2map_comp_id]])
apply (rule trans[OF arg_cong[OF F2.map_cong0]])
  apply (rule refl)

```

```

    apply (rule trans[OF o_apply])
    apply (rule id_apply)
    apply (rule trans[OF o_apply])
    apply (rule id_apply)
    apply (rule sym[OF o_apply])
done

lemmas IF1map_simps = o_eq_dest[OF IF1map]
lemmas IF2map_simps = o_eq_dest[OF IF2map]

lemma IFmap_unique:
  [[u o ctor1 = ctor1 o F1map f u v; v o ctor2 = ctor2 o F2map f u v]] ==>
    u = IF1map f ^ v = IF2map f
  apply (rule fold_unique)
  unfolding o_assoc[symmetric] F1.map_comp0[symmetric] F2.map_comp0[symmetric] id_o o_id
  apply assumption
  apply assumption
done

theorem IF1map_id: IF1map id = id
  apply (rule sym)
  apply (rule conjunct1[OF IFmap_unique])
  apply (rule trans[OF id_o])
  apply (rule trans[OF sym[OF o_id]])
  apply (rule arg_cong[OF sym[OF F1.map_id0]])
  apply (rule trans[OF id_o])
  apply (rule trans[OF sym[OF o_id]])
  apply (rule arg_cong[OF sym[OF F2.map_id0]])
done

theorem IF2map_id: IF2map id = id
  apply (rule sym)
  apply (rule conjunct2[OF IFmap_unique])
  apply (rule trans[OF id_o])
  apply (rule trans[OF sym[OF o_id]])
  apply (rule arg_cong[OF sym[OF F1.map_id0]])
  apply (rule trans[OF id_o])
  apply (rule trans[OF sym[OF o_id]])
  apply (rule arg_cong[OF sym[OF F2.map_id0]])
done

theorem IF1map_comp: IF1map (g o f) = IF1map g o IF1map f
  apply (rule sym)
  apply (rule conjunct1[OF IFmap_unique])
  apply (rule ext)
  apply (rule trans[OF o_apply])
  apply (rule trans[OF o_apply])
  apply (rule trans[OF arg_cong[OF IF1map_simps]])
  apply (rule trans[OF IF1map_simps])
  apply (rule trans[OF arg_cong[OF F1.map_comp]])
  apply (rule sym[OF o_apply])
  apply (rule ext)
  apply (rule trans[OF o_apply])
  apply (rule trans[OF o_apply])
  apply (rule trans[OF arg_cong[OF IF2map_simps]])
  apply (rule trans[OF IF2map_simps])
  apply (rule trans[OF arg_cong[OF F2.map_comp]])
  apply (rule sym[OF o_apply])
done

theorem IF2map_comp: IF2map (g o f) = IF2map g o IF2map f
  apply (rule sym)
  apply (tactic << rtac @{context} (Thm.permute_prem0 1 @{thm conjunct2[OF IFmap_unique]}) 1 >>)

```

```

apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule trans[OF o_apply])
apply (rule trans[OF arg_cong[OF IF2map_simps]])
apply (rule trans[OF IF2map_simps])
apply (rule trans[OF arg_cong[OF F2.map_comp]])
apply (rule sym[OF o_apply])
apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule trans[OF o_apply])
apply (rule trans[OF arg_cong[OF IF1map_simps]])
apply (rule trans[OF IF1map_simps])
apply (rule trans[OF arg_cong[OF F1.map_comp]])
apply (rule sym[OF o_apply])
done

```

The bound

abbreviation *IFbd* **where** $IFbd \equiv F1bd' + c F2bd'$

theorem *IFbd_card_order*: *card_order IFbd*

```

apply (rule card_order_csum)
apply (rule F1bd'_card_order)
apply (rule F2bd'_card_order)
done

```

lemma *IFbd_Cinfinite*: *Cinfinite IFbd*

```

apply (rule Cinfinite_csum1)
apply (rule F1bd'_Cinfinite)
done

```

lemmas *IFbd_cinfinite = conjunct1[OF IFbd_Cinfinite]*

The set operator

abbreviation *IF1col* **where** $IF1col \equiv (\lambda X. F1set1 X \cup (\bigcup F1set2 X \cup \bigcup F1set3 X))$

abbreviation *IF2col* **where** $IF2col \equiv (\lambda X. F2set1 X \cup (\bigcup F2set2 X \cup \bigcup F2set3 X))$

abbreviation *IF1set* **where** $IF1set \equiv fold1 IF1col IF2col$

abbreviation *IF2set* **where** $IF2set \equiv fold2 IF1col IF2col$

abbreviation *IF1in* **where** $IF1in A \equiv \{x. IF1set x \subseteq A\}$

abbreviation *IF2in* **where** $IF2in A \equiv \{x. IF2set x \subseteq A\}$

lemma *IF1set*: $IF1set o ctor1 = IF1col o (F1map id IF1set IF2set)$

```

apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule trans[OF fold1])
apply (rule sym[OF o_apply])
done

```

lemma *IF2set*: $IF2set o ctor2 = IF2col o (F2map id IF1set IF2set)$

```

apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule trans[OF fold2])
apply (rule sym[OF o_apply])
done

```

theorem *IF1set_simps*:

$IF1set (ctor1 x) = F1set1 x \cup ((\bigcup a \in F1set2 x. IF1set a) \cup (\bigcup a \in F1set3 x. IF2set a))$

```

apply (rule trans[OF o_eq_dest[OF IF1set]])
apply (rule arg_cong2[of _ _ _ (U)])
apply (rule trans[OF F1.set_map(1) trans[OF fun_cong[OF image_id] id_apply]])
apply (rule arg_cong2[of _ _ _ (U)])
apply (rule arg_cong[OF F1.set_map(2)])

```

```

apply (rule arg_cong[OF F1.set_map(3)])
done

```

theorem *IF2set_simps*:

```

 $IF2set\ (ctor2\ x) = F2set1\ x \cup ((\bigcup a \in F2set2\ x.\ IF1set\ a) \cup (\bigcup a \in F2set3\ x.\ IF2set\ a))$ 
apply (rule trans[OF o_eq_dest[OF IF2set]])
apply (rule arg_cong2[of _ _ _ _ (U)])
apply (rule trans[OF F2.set_map(1) trans[OF fun_cong[OF image_id] id_apply]])
apply (rule arg_cong2[of _ _ _ _ (U)])
apply (rule arg_cong[OF F2.set_map(2)])
apply (rule arg_cong[OF F2.set_map(3)])
done

```

lemmas *F1set1_IF1set = xt1(3)[OF IF1set_simps Un_upper1]*

lemmas *F1set2_IF1set = subset_trans[OF UN_upper subset_trans[OF Un_upper1 xt1(3)[OF IF1set_simps Un_upper2]]]*

lemmas *F1set3_IF1set = subset_trans[OF UN_upper subset_trans[OF Un_upper2 xt1(3)[OF IF1set_simps Un_upper2]]]*

lemmas *F2set1_IF2set = xt1(3)[OF IF2set_simps Un_upper1]*

lemmas *F2set2_IF2set = subset_trans[OF UN_upper subset_trans[OF Un_upper1 xt1(3)[OF IF2set_simps Un_upper2]]]*

lemmas *F2set3_IF2set = subset_trans[OF UN_upper subset_trans[OF Un_upper2 xt1(3)[OF IF2set_simps Un_upper2]]]*

The BNF conditions for IF

lemma *IFset_natural*:

```

 $f' (IF1set\ x) = IF1set\ (IF1map\ f\ x) \wedge f' (IF2set\ y) = IF2set\ (IF2map\ f\ y)$ 
apply (rule ctor_induct[of _ _ x y])

```

```

apply (rule trans)
apply (rule image_cong)
apply (rule IF1set_simps)
apply (rule refl)
apply (rule sym)
apply (rule trans[OF arg_cong[of _ _ IF1set, OF IF1map_simps] trans[OF IF1set_simps]])

```

```

apply (rule sym)
apply (rule trans)
apply (rule image_UN)
apply (rule arg_cong2[of _ _ _ _ (U)])
apply (rule sym)
apply (rule F1.set_map(1))

```

```

apply (rule trans)
apply (rule image_UN)
apply (rule arg_cong2[of _ _ _ _ (U)])
apply (rule trans)
apply (rule image_UN)
apply (rule trans)
apply (rule SUP_cong)
apply (rule refl)
apply (tactic << Goal.assume_rule_tac @ {context} 1 >>)
apply (rule sym)
apply (rule trans)
apply (rule SUP_cong)
apply (rule F1.set_map(2))
apply (rule refl)
apply (rule UN_simps(10))

```

```

apply (rule trans)
apply (rule image_UN)
apply (rule trans)
apply (rule SUP_cong)
apply (rule refl)
apply (tactic << Goal.assume_rule_tac @ {context} 1 >>)
apply (rule sym)

```

```

apply (rule trans)
apply (rule SUP_cong)
  apply (rule F1.set_map(3))
apply (rule refl)
apply (rule UN_simps(10))

```

```

apply (rule trans)
apply (rule image_cong)
  apply (rule IF2set_simps)
apply (rule refl)
apply (rule sym)
apply (rule trans[OF arg_cong[of _ _ IF2set, OF IF2map_simps] trans[OF IF2set_simps]])

```

```

apply (rule sym)
apply (rule trans)
  apply (rule image_Un)
apply (rule arg_cong2[of _ _ _ _ (U)])
apply (rule sym)
apply (rule F2.set_map(1))

```

```

apply (rule trans)
  apply (rule image_Un)
apply (rule arg_cong2[of _ _ _ _ (U)])

```

```

apply (rule trans)
  apply (rule image_UN)
apply (rule trans)
  apply (rule SUP_cong)
  apply (rule refl)
  apply (tactic << Goal.assume_rule_tac @{context} 1 >>)
apply (rule sym)
apply (rule trans)
  apply (rule SUP_cong)
  apply (rule F2.set_map(2))
apply (rule refl)
apply (rule UN_simps(10))

```

```

apply (rule trans)
  apply (rule image_UN)
apply (rule trans)
  apply (rule SUP_cong)
  apply (rule refl)
  apply (tactic << Goal.assume_rule_tac @{context} 1 >>)
apply (rule sym)
apply (rule trans)
  apply (rule SUP_cong)
  apply (rule F2.set_map(3))
apply (rule refl)
apply (rule UN_simps(10))
done

```

theorem *IF1set_natural*: $IF1set\ o\ (IF1map\ f) = image\ f\ o\ IF1set$

```

apply (rule ext)
apply (rule trans)
  apply (rule o_apply)
apply (rule sym)
apply (rule trans)
  apply (rule o_apply)
apply (rule conjunct1)
apply (rule IFset_natural)
done

```


theorem *IF2set_natural*: $IF2set \circ (IF2map \ f) = image \ f \circ IF2set$

```

apply (rule ext)
apply (rule trans)
  apply (rule o_apply)
apply (rule sym)
apply (rule trans)
  apply (rule o_apply)
apply (rule conjunct2)
apply (rule IFset_natural)
done

```

lemma *IFmap_cong*:

```

 $((\forall a \in IF1set \ x. f \ a = g \ a) \longrightarrow IF1map \ f \ x = IF1map \ g \ x) \wedge$ 
 $((\forall a \in IF2set \ y. f \ a = g \ a) \longrightarrow IF2map \ f \ y = IF2map \ g \ y)$ 
apply (rule ctor_induct[of _ _ x y])

```

```

apply (rule impI)
apply (rule trans)
  apply (rule IF1map_simps)
apply (rule trans)
  apply (rule arg_cong[OF F1.map_cong0])
    apply (erule bspec)
    apply (erule set_rev_mp)
    apply (rule F1set1_IF1set)
  apply (rule mp)
    apply (tactic << Goal.assume_rule_tac @{context} 1 >>)
  apply (rule ballI)
  apply (erule bspec)
  apply (erule set_rev_mp)
  apply (erule F1set2_IF1set)
apply (rule mp)
  apply (tactic << Goal.assume_rule_tac @{context} 1 >>)
apply (rule ballI)
apply (erule bspec)
apply (erule set_rev_mp)
apply (erule F1set3_IF1set)
apply (rule sym)
apply (rule IF1map_simps)

```

```

apply (rule impI)
apply (rule trans)
  apply (rule IF2map_simps)
apply (rule trans)
  apply (rule arg_cong[OF F2.map_cong0])
    apply (erule bspec)
    apply (erule set_rev_mp)
    apply (rule F2set1_IF2set)
  apply (rule mp)
    apply (tactic << Goal.assume_rule_tac @{context} 1 >>)
  apply (rule ballI)
  apply (erule bspec)
  apply (erule set_rev_mp)
  apply (erule F2set2_IF2set)
apply (rule mp)
  apply (tactic << Goal.assume_rule_tac @{context} 1 >>)
apply (rule ballI)
apply (erule bspec)
apply (erule set_rev_mp)
apply (erule F2set3_IF2set)
apply (rule sym)
apply (rule IF2map_simps)
done

```

theorem *IF1map_cong*:
 $(\bigwedge a. a \in IF1set\ x \implies f\ a = g\ a) \implies IF1map\ f\ x = IF1map\ g\ x$
apply (rule *mp*)
apply (rule *conjunct1*)
apply (rule *IFmap_cong*)
apply (rule *ballI*)
apply (tactic \ll *Goal.assume_rule_tac* @{*context*} 1 \gg)
done

theorem *IF2map_cong*:
 $(\bigwedge a. a \in IF2set\ x \implies f\ a = g\ a) \implies IF2map\ f\ x = IF2map\ g\ x$
apply (rule *mp*)
apply (rule *conjunct2*)
apply (rule *IFmap_cong*)
apply (rule *ballI*)
apply (tactic \ll *Goal.assume_rule_tac* @{*context*} 1 \gg)
done

lemma *IFset_bd*:
 $|IF1set\ (x :: 'a\ IF1)| \leq o\ IFbd \wedge |IF2set\ (y :: 'a\ IF2)| \leq o\ IFbd$
apply (rule *ctor_induct*[of _ _ x y])

apply (rule *ordIso_ordLeq_trans*)
apply (rule *card_of_ordIso_subst*)
apply (rule *IF1set_simps*)
apply (rule *Un_Cinfinite_bound*)
apply (rule *F1set1_bd'*)
apply (rule *Un_Cinfinite_bound*)
apply (rule *UNION_Cinfinite_bound*)
apply (rule *F1set2_bd'*)
apply (rule *ballI*)
apply (tactic \ll *Goal.assume_rule_tac* @{*context*} 1 \gg)
apply (rule *IFbd_Cinfinite*)
apply (rule *UNION_Cinfinite_bound*)
apply (rule *F1set3_bd'*)
apply (rule *ballI*)
apply (tactic \ll *Goal.assume_rule_tac* @{*context*} 1 \gg)
apply (rule *IFbd_Cinfinite*)
apply (rule *IFbd_Cinfinite*)
apply (rule *IFbd_Cinfinite*)

apply (rule *ordIso_ordLeq_trans*)
apply (rule *card_of_ordIso_subst*)
apply (rule *IF2set_simps*)
apply (rule *Un_Cinfinite_bound*)
apply (rule *F2set1_bd'*)
apply (rule *Un_Cinfinite_bound*)
apply (rule *UNION_Cinfinite_bound*)
apply (rule *F2set2_bd'*)
apply (rule *ballI*)
apply (tactic \ll *Goal.assume_rule_tac* @{*context*} 1 \gg)
apply (rule *IFbd_Cinfinite*)
apply (rule *UNION_Cinfinite_bound*)
apply (rule *F2set3_bd'*)
apply (rule *ballI*)
apply (tactic \ll *Goal.assume_rule_tac* @{*context*} 1 \gg)
apply (rule *IFbd_Cinfinite*)
apply (rule *IFbd_Cinfinite*)
apply (rule *IFbd_Cinfinite*)
done

lemmas *IF1set_bd = conjunct1[OF IFset_bd]*
lemmas *IF2set_bd = conjunct2[OF IFset_bd]*

definition *IF1rel* where

```
IF1rel R =  
  (BNF_Def.Grp (IF1in (Collect (case_prod R))) (IF1map fst)) ^--1 OO  
  (BNF_Def.Grp (IF1in (Collect (case_prod R))) (IF1map snd))
```

definition *IF2rel* where

```
IF2rel R =  
  (BNF_Def.Grp (IF2in (Collect (case_prod R))) (IF2map fst)) ^--1 OO  
  (BNF_Def.Grp (IF2in (Collect (case_prod R))) (IF2map snd))
```

lemma *in_IF1rel*:

```
IF1rel R x y  $\longleftrightarrow$  ( $\exists z. z \in \text{IF1in (Collect (case\_prod R))} \wedge \text{IF1map fst } z = x \wedge \text{IF1map snd } z = y$ )  
unfolding IF1rel_def by (rule predicate2_eqD[OF OO_Grp_alt])
```

lemma *in_IF2rel*:

```
IF2rel R x y  $\longleftrightarrow$  ( $\exists z. z \in \text{IF2in (Collect (case\_prod R))} \wedge \text{IF2map fst } z = x \wedge \text{IF2map snd } z = y$ )  
unfolding IF2rel_def by (rule predicate2_eqD[OF OO_Grp_alt])
```

lemma *IF1rel_F1rel*: *IF1rel* R (ctor1 a) (ctor1 b) \longleftrightarrow *F1rel* R (*IF1rel* R) (*IF2rel* R) a b

```
apply (rule iffI)  
apply (tactic << dtac @{context} (@{thm in_IF1rel[THEN iffD1]} 1 >>)+  
apply (erule exE conjE CollectE)+  
apply (rule iffD2)  
apply (rule F1.in_rel)  
apply (rule exI)  
apply (rule conjI)  
apply (rule CollectI)  
apply (rule conjI)  
apply (rule ord_eq_le_trans)  
  apply (rule F1.set_map(1))  
  apply (rule ord_eq_le_trans)  
  apply (rule trans[OF fun_cong[OF image_id] id_apply])  
  apply (rule subset_trans)  
  apply (rule F1set1_IF1set)  
  apply (erule ord_eq_le_trans[OF arg_cong[OF ctor1_dtor1]])
```

```
apply (rule conjI)  
  apply (rule ord_eq_le_trans)  
    apply (rule F1.set_map(2))  
  apply (rule image_subsetI)  
  apply (rule CollectI)  
  apply (rule case_prodI)  
  apply (rule iffD2)  
  apply (rule in_IF1rel)  
  apply (rule exI)  
  apply (rule conjI)  
  apply (rule CollectI)  
  apply (erule subset_trans[OF F1set2_IF1set])  
  apply (erule ord_eq_le_trans[OF arg_cong[OF ctor1_dtor1]])  
  apply (rule conjI)  
  apply (rule refl)  
  apply (rule refl)
```

```
apply (rule ord_eq_le_trans)  
  apply (rule F1.set_map(3))  
  apply (rule image_subsetI)  
  apply (rule CollectI)  
  apply (rule case_prodI)  
  apply (rule iffD2)  
  apply (rule in_IF2rel)  
  apply (rule exI)  
  apply (rule conjI)
```

```

apply (rule CollectI)
apply (rule subset_trans)
apply (rule F1set3_IF1set)
apply assumption
apply (erule ord_eq_le_trans[OF arg_cong[OF ctor1_dtor1]])
apply (rule conjI)
apply (rule refl)
apply (rule refl)
apply (rule conjI)

```

```

apply (rule trans)
apply (rule F1.map_comp)
apply (rule trans)
apply (rule F1.map_cong0)
apply (rule fun_cong[OF o_id])
apply (rule trans)
apply (rule o_apply)
apply (rule fst_conv)
apply (rule trans)
apply (rule o_apply)
apply (rule fst_conv)
apply (rule iffD1[OF ctor1_diff])
apply (rule trans)
apply (rule sym)
apply (rule IF1map_simps)
apply (erule trans[OF arg_cong[OF ctor1_dtor1]])

```

```

apply (rule trans)
apply (rule F1.map_comp)
apply (rule trans)
apply (rule F1.map_cong0)
apply (rule fun_cong[OF o_id])
apply (rule trans)
apply (rule o_apply)
apply (rule snd_conv)
apply (rule trans)
apply (rule o_apply)
apply (rule snd_conv)
apply (rule iffD1[OF ctor1_diff])
apply (rule trans)
apply (rule sym)
apply (rule IF1map_simps)
apply (erule trans[OF arg_cong[OF ctor1_dtor1]])

```

```

apply (tactic << dtac @{context} (@{thm F1.in_rel[THEN iffD1]}) 1 >>)
apply (erule exE conjE CollectE)+
apply (rule iffD2)
apply (rule in_IF1rel)
apply (rule exI)
apply (rule conjI)
apply (rule CollectI)
apply (rule ord_eq_le_trans)
apply (rule IF1set_simps)
apply (rule Un_least)
apply (rule ord_eq_le_trans)
apply (rule box_equals[OF _ refl])
apply (rule F1.set_map(1))
apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply assumption
apply (rule Un_least)
apply (rule ord_eq_le_trans)
apply (rule SUP_cong[OF _ refl])

```

```

  apply (rule F1.set_map(2))
  apply (rule UN_least)
  apply (drule set_rev_mp)
  apply (erule image_mono)
  apply (erule imageE)
  apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
  apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
  apply hypsubst
  apply (tactic « dtac @{context} (@{thm in_IF1rel[THEN iffD1]}) 1 »)
  apply (drule someI_ex)
  apply (erule conjE)+
  apply (erule CollectD)

  apply (rule ord_eq_le_trans)
  apply (rule SUP_cong[OF _ refl])
  apply (rule F1.set_map(3))
  apply (rule UN_least)
  apply (drule set_rev_mp)
  apply (erule image_mono)
  apply (erule imageE)
  apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
  apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
  apply hypsubst
  apply (tactic « dtac @{context} (@{thm in_IF2rel[THEN iffD1]}) 1 »)
  apply (drule someI_ex)
  apply (erule conjE)+
  apply (erule CollectD)

  apply (rule conjI)
  apply (rule trans)
  apply (rule IF1map_simps)
  apply (rule iffD2[OF ctor1_diff])
  apply (rule trans)
  apply (rule F1.map_comp)
  apply (rule trans)
  apply (rule F1.map_cong0)
  apply (rule fun_cong[OF o_id])
  apply (rule trans[OF o_apply])
  apply (drule set_rev_mp)
  apply assumption
  apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
  apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
  apply hypsubst
  apply (tactic « dtac @{context} (@{thm in_IF1rel[THEN iffD1]}) 1 »)
  apply (drule someI_ex)
  apply (erule conjE)+
  apply assumption
  apply (rule trans[OF o_apply])
  apply (drule set_rev_mp)
  apply assumption
  apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
  apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
  apply hypsubst
  apply (tactic « dtac @{context} (@{thm in_IF2rel[THEN iffD1]}) 1 »)
  apply (drule someI_ex)
  apply (erule conjE)+
  apply assumption
  apply assumption

  apply (rule trans)
  apply (rule IF1map_simps)
  apply (rule iffD2[OF ctor1_diff])
  apply (rule trans)

```

```

apply (rule F1.map_comp)
apply (rule trans)
apply (rule F1.map_cong0)
  apply (rule fun_cong[OF o_id])
  apply (rule trans[OF o_apply])
apply (drule set_rev_mp)
  apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format conjE])+
apply hypsubst
apply (tactic << dtac @{context} (@{thm in_IF1rel[THEN iffD1]} 1 >>))
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply (rule trans[OF o_apply])
apply (drule set_rev_mp)
  apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format conjE])+
apply hypsubst
apply (tactic << dtac @{context} (@{thm in_IF2rel[THEN iffD1]} 1 >>))
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply assumption
done

```

lemma *IF2rel_F2rel*: $IF2rel\ R\ (ctor2\ a)\ (ctor2\ b) \longleftrightarrow F2rel\ R\ (IF1rel\ R)\ (IF2rel\ R)\ a\ b$

```

apply (rule iffI)
apply (tactic << dtac @{context} (@{thm in_IF2rel[THEN iffD1]} 1 >>))+
apply (erule exE conjE CollectE)+
apply (rule iffD2)
apply (rule F2.in_rel)
apply (rule exI)
apply (rule conjI)
apply (rule CollectI)
apply (rule conjI)
  apply (rule ord_eq_le_trans)
    apply (rule F2.set_map(1))
  apply (rule ord_eq_le_trans)
    apply (rule trans[OF fun_cong[OF image_id id_apply]])
  apply (rule subset_trans)
    apply (rule F2set1_IF2set)
  apply (erule ord_eq_le_trans[OF arg_cong[OF ctor2_dtor2]])

apply (rule conjI)
  apply (rule ord_eq_le_trans)
    apply (rule F2.set_map(2))
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule case_prodI)
  apply (rule iffD2)
  apply (rule in_IF1rel)
  apply (rule exI)
  apply (rule conjI)
  apply (rule CollectI)
  apply (rule subset_trans)
    apply (rule F2set2_IF2set)
  apply assumption
  apply (erule ord_eq_le_trans[OF arg_cong[OF ctor2_dtor2]])
apply (rule conjI)
  apply (rule refl)
apply (rule refl)

```

```

apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(3))
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule case_prodI)
apply (rule iffD2)
  apply (rule in_IF2rel)
apply (rule exI)
apply (rule conjI)
  apply (rule CollectI)
  apply (rule subset_trans)
  apply (rule F2set3_IF2set)
  apply assumption
apply (erule ord_eq_le_trans[OF arg_cong[OF ctor2_dtor2]])
apply (rule conjI)
  apply (rule refl)
apply (rule refl)
apply (rule conjI)

```

```

apply (rule trans)
  apply (rule F2.map_comp)
apply (rule trans)
  apply (rule F2.map_cong0)
    apply (rule fun_cong[OF o_id])
  apply (rule trans)
    apply (rule o_apply)
  apply (rule fst_conv)
apply (rule trans)
  apply (rule o_apply)
apply (rule fst_conv)
apply (rule iffD1[OF ctor2_diff])
apply (rule trans)
  apply (rule sym)
  apply (rule IF2map_simps)
apply (erule trans[OF arg_cong[OF ctor2_dtor2]])

```

```

apply (rule trans)
apply (rule F2.map_comp)
apply (rule trans)
apply (rule F2.map_cong0)
  apply (rule fun_cong[OF o_id])
  apply (rule trans)
    apply (rule o_apply)
  apply (rule snd_conv)
apply (rule trans)
  apply (rule o_apply)
apply (rule snd_conv)
apply (rule iffD1[OF ctor2_diff])
apply (rule trans)
  apply (rule sym)
  apply (rule IF2map_simps)
apply (erule trans[OF arg_cong[OF ctor2_dtor2]])

```

```

apply (tactic << dtac @{context} (@{thm F2.in_rel[THEN iffD1]} 1 >>))
apply (erule exE conjE CollectE)+
apply (rule iffD2)
  apply (rule in_IF2rel)
apply (rule exI)
apply (rule conjI)
apply (rule CollectI)
apply (rule ord_eq_le_trans)

```

```

apply (rule IF2set_simps)
apply (rule Un_least)
apply (rule ord_eq_le_trans)
  apply (rule trans)
    apply (rule trans)
      apply (rule arg_cong[OF dtor2_ctor2])
      apply (rule F2.set_map(1))
    apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply assumption
apply (rule Un_least)
apply (rule ord_eq_le_trans)
  apply (rule trans[OF arg_cong[OF dtor2_ctor2]])
  apply (rule arg_cong[OF F2.set_map(2)])
apply (rule UN_least)
apply (drule set_rev_mp)
  apply (erule image_mono)
apply (erule imageE)
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_proxE iffD1[OF prod.inject, elim_format] conjE)+
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (tactic << dtac @ {context} (@ {thm in_IF1rel[THEN iffD1]}) 1 >>)
apply (drule someI_ex)
apply (erule conjE)+
apply (erule CollectD)

apply (rule ord_eq_le_trans)
apply (rule trans[OF arg_cong[OF dtor2_ctor2]])
apply (rule arg_cong[OF F2.set_map(3)])
apply (rule UN_least)
apply (drule set_rev_mp)
apply (erule image_mono)
apply (erule imageE)
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_proxE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (tactic << dtac @ {context} (@ {thm in_IF2rel[THEN iffD1]}) 1 >>)
apply (drule someI_ex)
apply (erule exE conjE)+
apply (erule CollectD)

apply (rule conjI)
apply (rule trans)
  apply (rule arg_cong[OF dtor2_ctor2])
apply (rule trans)
  apply (rule IF2map_simps)
apply (rule iffD2)
apply (rule ctor2_diff)
apply (rule trans)
  apply (rule F2.map_comp)
apply (rule trans)
apply (rule F2.map_cong0)
  apply (rule fun_cong[OF o_id])
  apply (rule trans[OF o_apply])
  apply (drule set_rev_mp)
  apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_proxE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (tactic << dtac @ {context} (@ {thm in_IF1rel[THEN iffD1]}) 1 >>)
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply (rule trans[OF o_apply])

```



```

apply (drule set_rev_mp)
apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (tactic << dtac @{context} (@{thm in_IF2rel[THEN iffD1]}) 1 >>)
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply assumption

```

```

apply (rule trans)
apply (rule arg_cong[OF dtor2_ctor2])
apply (rule trans)
apply (rule IF2map_simps)
apply (rule iffD2)
apply (rule ctor2_diff)
apply (rule trans)
apply (rule F2.map_comp)
apply (rule trans)
apply (rule F2.map_cong0)
apply (rule fun_cong[OF o_id])
apply (rule trans[OF o_apply])
apply (drule set_rev_mp)
apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (tactic << dtac @{context} (@{thm in_IF1rel[THEN iffD1]}) 1 >>)
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply (rule trans[OF o_apply])
apply (drule set_rev_mp)
apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (tactic << dtac @{context} (@{thm in_IF2rel[THEN iffD1]}) 1 >>)
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply assumption
done

```

lemma Irel_induct:

```

assumes IH1:  $\forall x y. F1rel\ P1\ P2\ P3\ x\ y \longrightarrow P2\ (ctor1\ x)\ (ctor1\ y)$ 
and IH2:  $\forall x y. F2rel\ P1\ P2\ P3\ x\ y \longrightarrow P3\ (ctor2\ x)\ (ctor2\ y)$ 
shows  $IF1rel\ P1 \leq P2 \wedge IF2rel\ P1 \leq P3$ 
unfolding le_fun_def le_bool_def all_simps(1,2)[symmetric]
apply (rule allI)+
apply (rule ctor_induct2)
apply (rule impI)
apply (drule iffD1[OF IF1rel_F1rel])
apply (rule mp[OF spec2[OF IH1]])
apply (erule F1.rel_mono_strong0)
apply (rule ballI[OF ballI[OF imp_refl]])
apply (drule asm_rl)
apply (erule thin_rl)
apply (rule ballI[OF ballI])
apply assumption
apply (erule thin_rl)
apply (drule asm_rl)

```

apply (rule ballI[OF ballI])
apply assumption

apply (rule impI)
apply (drule iffD1[OF IF2rel_F2rel])
apply (rule mp[OF spec2[OF IH2]])
apply (erule F2.rel_mono_strong0)
apply (rule ballI[OF ballI[OF imp_refl]])
apply (drule asm_rl)
apply (erule thin_rl)
apply (rule ballI[OF ballI])
apply assumption
apply (erule thin_rl)
apply (drule asm_rl)
apply (rule ballI[OF ballI])
apply assumption
done

lemma le_IFrel_Comp:

$((IF1rel\ R\ OO\ IF1rel\ S)\ x1\ y1 \longrightarrow IF1rel\ (R\ OO\ S)\ x1\ y1) \wedge$
 $((IF2rel\ R\ OO\ IF2rel\ S)\ x2\ y2 \longrightarrow IF2rel\ (R\ OO\ S)\ x2\ y2)$
apply (rule ctor_induct2[of _ _ x1 y1 x2 y2])
apply (rule impI)
apply (erule nchotomy_relcomppE[OF ctor1_nchotomy])
apply (drule iffD1[OF IF1rel_F1rel])
apply (drule iffD1[OF IF1rel_F1rel])
apply (rule iffD2[OF IF1rel_F1rel])
apply (rule F1.rel_mono_strong0)
apply (rule iffD2[OF predicate2_eqD[OF F1.rel_compp]])
apply (rule relcomppI)
apply assumption
apply assumption
apply (rule ballI impI)+
apply assumption
apply (rule ballI)+
apply assumption
apply (rule ballI)+
apply assumption

apply (rule impI)
apply (erule nchotomy_relcomppE[OF ctor2_nchotomy])
apply (drule iffD1[OF IF2rel_F2rel])
apply (drule iffD1[OF IF2rel_F2rel])
apply (rule iffD2[OF IF2rel_F2rel])
apply (rule F2.rel_mono_strong0)
apply (rule iffD2[OF predicate2_eqD[OF F2.rel_compp]])
apply (rule relcomppI)
apply assumption
apply assumption
apply (rule ballI impI)+
apply assumption
apply (rule ballI)+
apply assumption
apply (rule ballI)+
apply assumption
done

lemma le_IF1rel_Comp: $IF1rel\ R1\ OO\ IF1rel\ R2 \leq IF1rel\ (R1\ OO\ R2)$
by (rule predicate2I) (erule mp[OF conjunct1[OF le_IFrel_Comp]])

lemma le_IF2rel_Comp: $IF2rel\ R1\ OO\ IF2rel\ R2 \leq IF2rel\ (R1\ OO\ R2)$
by (rule predicate2I) (erule mp[OF conjunct2[OF le_IFrel_Comp]])

context includes *lifting_syntax*
begin

lemma *fold_transfer*:

```
((F1rel R S T ==> S) ==> (F2rel R S T ==> T) ==> IF1rel R ==> S) fold1 fold1 ^
((F1rel R S T ==> S) ==> (F2rel R S T ==> T) ==> IF2rel R ==> T) fold2 fold2
unfolding rel_fun_def_butlast all_conj_distrib[symmetric] imp_conjR[symmetric]
unfolding rel_fun_iff_leq_vimage2p
apply (rule allI_impI)+
apply (rule Irel_induct)
apply (rule allI_impI_vimage2pI)+
apply (unfold fold1 fold2) [1]
apply (erule predicate2D_vimage2p)
apply (rule rel_funD[OF rel_funD[OF rel_funD[OF F1.map_transfer]]])
  apply (rule id_transfer)
  apply (rule vimage2p_rel_fun)
  apply (rule vimage2p_rel_fun)
apply assumption
```

```
apply (rule allI_impI_vimage2pI)+
apply (unfold fold1 fold2) [1]
apply (erule predicate2D_vimage2p)
apply (rule rel_funD[OF rel_funD[OF rel_funD[OF F2.map_transfer]]])
  apply (rule id_transfer)
  apply (rule vimage2p_rel_fun)
  apply (rule vimage2p_rel_fun)
apply assumption
done
```

end

definition *IF1wit* x = *ctor1* (*wit2_F1* x (*ctor2* *wit_F2*))

definition *IF2wit* = *ctor2* *wit_F2*

lemma *IF1wit*: $x \in IF1set (IF1wit\ y) \implies x = y$

```
unfolding IF1wit_def
by (elim UnE F1.wit2[elim_format] F2.wit[elim_format] UN_E FalseE |
  rule refl | hypsubst | assumption | unfold IF1set_simps IF2set_simps)+
```

lemma *IF2wit*: $x \in IF2set\ IF2wit \implies False$

```
unfolding IF2wit_def
by (elim UnE F2.wit[elim_format] UN_E FalseE |
  rule refl | hypsubst | assumption | unfold IF2set_simps)+
```

ML <<

```
BNF_FP_Util.mk_xtor_co_iter_o_map_thms BNF_Util.Least_FP false 1 @ {thm fold_unique}
@ {thms IF1map IF2map} (map (BNF_Tactics.mk_pointfree2 @ {context}) @ {thms fold1 fold2})
@ {thms F1.map_comp0[symmetric] F2.map_comp0[symmetric]} @ {thms F1.map_cong0 F2.map_cong0}
>>
```

ML <<

```
BNF_FP_Util.mk_xtor_co_iter_o_map_thms BNF_Util.Least_FP true 1 @ {thm rec_unique}
@ {thms IF1map IF2map} (map (BNF_Tactics.mk_pointfree2 @ {context}) @ {thms rec1 rec2})
@ {thms F1.map_comp0[symmetric] F2.map_comp0[symmetric]} @ {thms F1.map_cong0 F2.map_cong0}
>>
```

bnf 'a *IF1*

```
map: IF1map
sets: IF1set
bd: IFbd
wits: IF1wit
rel: IF1rel
```

```

    apply -
    apply (rule IF1map_id)
    apply (rule IF1map_comp)
    apply (erule IF1map_cong)
    apply (rule IF1set_natural)
    apply (rule IFbd_card_order)
    apply (rule IFbd_cinfinite)
    apply (rule IF1set_bd)
    apply (rule le_IF1rel_Comp)
    apply (rule IF1rel_def[unfolded OO_Grp_alt mem_Collect_eq])
    apply (erule IF1wit)
done

bnf 'a IF2
map: IF2map
sets: IF2set
bd: IFbd
wits: IF2wit
rel: IF2rel
    apply -
    apply (rule IF2map_id)
    apply (rule IF2map_comp)
    apply (erule IF2map_cong)
    apply (rule IF2set_natural)
    apply (rule IFbd_card_order)
    apply (rule IFbd_cinfinite)
    apply (rule IF2set_bd)
    apply (rule le_IF2rel_Comp)
    apply (rule IF2rel_def[unfolded OO_Grp_alt mem_Collect_eq])
    apply (erule IF2wit)
done

```

2 Greatest Fixpoint (a.k.a. Codatatype)

```

'b1 = ('a, 'b1, 'b2) F1
'b2 = ('a, 'b1, 'b2) F2

```

To build a witness scenario, let us assume

```

('a, 'b1, 'b2) F1 = 'a * 'b1 + 'a * 'b2
('a, 'b1, 'b2) F2 = unit + 'b1 * 'b2

```

```

ML (open Ctr_Sugar_Util)
declare [[bnf_internals]]
bnf-axiomatization (F1set1: 'a, F1set2: 'b1, F1set3: 'b2) F1
[wits: 'a ⇒ 'b1 ⇒ ('a, 'b1, 'b2) F1 'a ⇒ 'b2 ⇒ ('a, 'b1, 'b2) F1]
for map: F1map rel: F1rel

```

```

bnf-axiomatization (F2set1: 'a, F2set2: 'b1, F2set3: 'b2) F2
[wits: ('a, 'b1, 'b2) F2]
for map: F2map rel: F2rel

```

```

lemma F1rel_cong: [[R1 = S1; R2 = S2; R3 = S3]] ⇒ F1rel R1 R2 R3 = F1rel S1 S2 S3
by hypsubst rule

```

```

lemma F2rel_cong: [[R1 = S1; R2 = S2; R3 = S3]] ⇒ F2rel R1 R2 R3 = F2rel S1 S2 S3
by hypsubst rule

```

```

abbreviation F1in :: 'a1 set ⇒ 'a2 set ⇒ 'a3 set ⇒ (('a1, 'a2, 'a3) F1) set where
F1in A1 A2 A3 ≡ {x. F1set1 x ⊆ A1 ∧ F1set2 x ⊆ A2 ∧ F1set3 x ⊆ A3}

```

```

abbreviation F2in :: 'a1 set ⇒ 'a2 set ⇒ 'a3 set ⇒ (('a1, 'a2, 'a3) F2) set where
F2in A1 A2 A3 ≡ {x. F2set1 x ⊆ A1 ∧ F2set2 x ⊆ A2 ∧ F2set3 x ⊆ A3}

```

lemma *F1map_comp_id*: $F1map\ g1\ g2\ g3\ (F1map\ id\ f2\ f3\ x) = F1map\ g1\ (g2\ o\ f2)\ (g3\ o\ f3)\ x$
apply (*rule trans*)
apply (*rule F1.map_comp*)
unfolding *o_id*
apply (*rule refl*)
done

lemmas *F1in_mono23* = *F1.in_mono*[*OF subset_refl*]
lemmas *F1in_mono23'* = *set_mp*[*OF F1in_mono23*]

lemma *F1map_congL*: $\llbracket \forall a \in F1set2\ x.\ f\ a = a; \forall a \in F1set3\ x.\ g\ a = a \rrbracket \implies$
 $F1map\ id\ f\ g\ x = x$
apply (*rule trans*)
apply (*rule F1.map_cong0*)
apply (*rule refl*)
apply (*rule trans*)
apply (*erule bspec*)
apply *assumption*
apply (*rule sym*)
apply (*rule id_apply*)
apply (*rule trans*)
apply (*erule bspec*)
apply *assumption*
apply (*rule sym*)
apply (*rule id_apply*)
apply (*rule F1.map_id*)
done

lemma *F2map_comp_id*: $F2map\ g1\ g2\ g3\ (F2map\ id\ f2\ f3\ x) = F2map\ g1\ (g2\ o\ f2)\ (g3\ o\ f3)\ x$
apply (*rule trans*)
apply (*rule F2.map_comp*)
unfolding *o_id*
apply (*rule refl*)
done

lemmas *F2in_mono23* = *F2.in_mono*[*OF subset_refl*]
lemmas *F2in_mono23'* = *set_mp*[*OF F2in_mono23*]

lemma *F2map_congL*: $\llbracket \forall a \in F2set2\ x.\ f\ a = a; \forall a \in F2set3\ x.\ g\ a = a \rrbracket \implies$
 $F2map\ id\ f\ g\ x = x$
apply (*rule trans*)
apply (*rule F2.map_cong0*)
apply (*rule refl*)
apply (*rule trans*)
apply (*erule bspec*)
apply *assumption*
apply (*rule sym*)
apply (*rule id_apply*)
apply (*rule trans*)
apply (*erule bspec*)
apply *assumption*
apply (*rule sym*)
apply (*rule id_apply*)
apply (*rule F2.map_id*)
done

2.1 Coalgebra

definition *coalg where*

$coalg\ B1\ B2\ s1\ s2 =$
 $(\forall a \in B1.\ s1\ a \in F1in\ (UNIV :: 'a\ set)\ B1\ B2) \wedge (\forall a \in B2.\ s2\ a \in F2in\ (UNIV :: 'a\ set)\ B1\ B2)$

lemmas *coalg_F1in* = *bspec*[*OF conjunct1*[*OF iffD1*[*OF coalg_def*]]]
lemmas *coalg_F2in* = *bspec*[*OF conjunct2*[*OF iffD1*[*OF coalg_def*]]]

```

lemma coalg_F1set2:
  [[coalg B1 B2 s1 s2; a ∈ B1]] ⇒ F1set2 (s1 a) ⊆ B1
  apply (tactic << dtac @{context} @{thm iffD1[OF coalg_def]} 1 >>)
  apply (erule conjE)
  apply (drule bspec[rotated])
  apply assumption
  apply (erule CollectE conjE)+
  apply assumption
  done

```

```

lemma coalg_F1set3:
  [[coalg B1 B2 s1 s2; a ∈ B1]] ⇒ F1set3 (s1 a) ⊆ B2
  apply (tactic << dtac @{context} @{thm iffD1[OF coalg_def]} 1 >>)
  apply (erule conjE)
  apply (drule bspec[rotated])
  apply assumption
  apply (erule CollectE conjE)+
  apply assumption
  done

```

```

lemma coalg_F2set2:
  [[coalg B1 B2 s1 s2; a ∈ B2]] ⇒ F2set2 (s2 a) ⊆ B1
  apply (tactic << dtac @{context} @{thm iffD1[OF coalg_def]} 1 >>)
  apply (erule conjE)
  apply (drule bspec[rotated])
  apply assumption
  apply (erule CollectE conjE)+
  apply assumption
  done

```

```

lemma coalg_F2set3:
  [[coalg B1 B2 s1 s2; a ∈ B2]] ⇒ F2set3 (s2 a) ⊆ B2
  apply (tactic << dtac @{context} @{thm iffD1[OF coalg_def]} 1 >>)
  apply (erule conjE)
  apply (drule bspec[rotated])
  apply assumption
  apply (erule CollectE conjE)+
  apply assumption
  done

```

2.2 Type-coalgebra

abbreviation *tcoalg s1 s2* ≡ *coalg UNIV UNIV s1 s2*

```

lemma tcoalg: tcoalg s1 s2
  apply (tactic << rtac @{context} (@{thm coalg_def} RS iffD2) 1 >>)
  apply (rule conjI)
  apply (rule ballI)
  apply (rule CollectI)
  apply (rule conjI)
  apply (rule subset_UNIV)
  apply (rule conjI)
  apply (rule subset_UNIV)
  apply (rule subset_UNIV)
  apply (rule ballI)
  apply (rule CollectI)
  apply (rule conjI)
  apply (rule subset_UNIV)
  apply (rule conjI)
  apply (rule subset_UNIV)
  apply (rule subset_UNIV)
  done

```

2.3 Morphism

definition *mor where*

```
mor B1 B2 s1 s2 B1' B2' s1' s2' f g =
  (((∀ a ∈ B1. f a ∈ B1') ∧ (∀ a ∈ B2. g a ∈ B2')) ∧
  ((∀ z ∈ B1. F1map (id :: 'a ⇒ 'a) f g (s1 z) = s1' (f z)) ∧
  (∀ z ∈ B2. F2map (id :: 'a ⇒ 'a) f g (s2 z) = s2' (g z))))
```

lemma *mor_image1*: $\text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g \Longrightarrow f \ ' \ B1 \subseteq B1'$

```
apply (tactic << dtac @{context} @{thm iffD1[OF mor_def]} 1 >>)
apply (erule conjE)+
apply (rule image_subsetI)
apply (erule bspec)
apply assumption
done
```

lemma *mor_image2*: $\text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g \Longrightarrow g \ ' \ B2 \subseteq B2'$

```
apply (tactic << dtac @{context} @{thm iffD1[OF mor_def]} 1 >>)
apply (erule conjE)+
apply (rule image_subsetI)
apply (erule bspec)
apply assumption
done
```

lemmas *mor_image1'* = $\text{set_mp}[OF \ \text{mor_image1} \ \text{imageI}]$
lemmas *mor_image2'* = $\text{set_mp}[OF \ \text{mor_image2} \ \text{imageI}]$

lemma *morE1*: $\llbracket \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g; z \in B1 \rrbracket$

```
⇒ F1map id f g (s1 z) = s1' (f z)
apply (tactic << dtac @{context} @{thm iffD1[OF mor_def]} 1 >>)
apply (erule conjE)+
apply (erule bspec)
apply assumption
done
```

lemma *morE2*: $\llbracket \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f \ g; z \in B2 \rrbracket$

```
⇒ F2map id f g (s2 z) = s2' (g z)
apply (tactic << dtac @{context} @{thm iffD1[OF mor_def]} 1 >>)
apply (erule conjE)+
apply (erule bspec)
apply assumption
done
```

lemma *mor_incl*: $\llbracket B1 \subseteq B1'; B2 \subseteq B2' \rrbracket \Longrightarrow \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1 \ s2 \ \text{id} \ \text{id}$

```
apply (tactic << rtac @{context} (@{thm mor_def} RS iffD2) 1 >>)
apply (rule conjI)
apply (rule conjI)
apply (rule ballI)
apply (rule ssubst_mem[OF id_apply])
apply (erule set_mp)
apply assumption
```

```
apply (rule ballI)
apply (rule ssubst_mem[OF id_apply])
apply (erule set_mp)
apply assumption
```

```
apply (rule conjI)
apply (rule ballI)
apply (rule trans[OF F1.map_id])
apply (rule sym)
apply (rule arg_cong[OF id_apply])
apply (rule ballI)
apply (rule trans[OF F2.map_id])
```

```

apply (rule sym)
apply (rule arg_cong[OF id_apply])
done

```

```

lemmas mor_id = mor_incl[OF subset_refl subset_refl]

```

```

lemma mor_comp:

```

```

  [[mor B1 B2 s1 s2 B1' B2' s1' s2' f g;
   mor B1' B2' s1' s2' B1'' B2'' s1'' s2'' f' g'] ==>
   mor B1 B2 s1 s2 B1'' B2'' s1'' s2'' (f' o f) (g' o g)]
apply (tactic << rtac @{context} (@{thm mor_def} RS iffD2) 1 >>))
apply (rule conjI)

```

```

apply (rule conjI)
apply (rule ballI)
apply (rule ssubst_mem[OF o_apply])
apply (erule mor_image1')
apply (erule mor_image1')
apply assumption

```

```

apply (rule ballI)
apply (rule ssubst_mem[OF o_apply])
apply (erule mor_image2')
apply (erule mor_image2')
apply assumption

```

```

apply (rule conjI)
apply (rule ballI)
apply (tactic << stac @{context} @{thm o_apply} 1 >>))
apply (rule trans)
apply (rule sym[OF F1map_comp_id])
apply (rule trans)
apply (erule arg_cong[OF morE1])
apply assumption
apply (erule morE1)
apply (erule mor_image1')
apply assumption

```

```

apply (rule ballI)
apply (tactic << stac @{context} @{thm o_apply} 1 >>))
apply (rule trans)
apply (rule sym[OF F2map_comp_id])
apply (rule trans)
apply (erule arg_cong[OF morE2])
apply assumption
apply (erule morE2)
apply (erule mor_image2')
apply assumption
done

```

```

lemma mor_cong: [[ f' = f; g' = g; mor B1 B2 s1 s2 B1' B2' s1' s2' f g ] ==>
  mor B1 B2 s1 s2 B1' B2' s1' s2' f' g' ]
apply (tactic << hyp_subst_tac @{context} 1 >>))
apply assumption
done

```

```

lemma mor_UNIV: mor UNIV UNIV s1 s2 UNIV UNIV s1' s2' f1 f2 <=>
  F1map id f1 f2 o s1 = s1' o f1 & F2map id f1 f2 o s2 = s2' o f2
apply (rule iffI)
apply (rule conjI)
apply (rule ext)
apply (rule trans)
apply (rule trans)

```



```

  apply (rule o_apply)
  apply (erule morE1)
  apply (rule UNIV_I)
  apply (rule sym[OF o_apply])
  apply (rule ext)
  apply (rule trans)
  apply (rule trans)
  apply (rule o_apply)
  apply (erule morE2)
  apply (rule UNIV_I)
  apply (rule sym[OF o_apply])

  apply (tactic << rtac @context (thm mor_def) RS iffD2) 1 >>)
  apply (rule conjI)
  apply (rule conjI)
  apply (rule ballI)
  apply (rule UNIV_I)
  apply (rule ballI)
  apply (rule UNIV_I)
  apply (rule conjI)
  apply (tactic << dtac @context (BNF_Util.mk_conjunctN 2 1) 1 >>)
  apply (rule ballI)
  apply (erule o_eq_dest)
  apply (tactic << dtac @context (BNF_Util.mk_conjunctN 2 2) 1 >>)
  apply (rule ballI)
  apply (erule o_eq_dest)
done

```

lemma mor_str:

```

  mor UNIV UNIV s1 s2 UNIV UNIV (F1map id s1 s2) (F2map id s1 s2) s1 s2
  apply (rule iffD2)
  apply (rule mor_UNIV)
  apply (rule conjI)
  apply (rule refl)
  apply (rule refl)
done

```

lemma mor_case_sum:

```

  mor UNIV UNIV s1 s2 UNIV UNIV (case_sum (F1map id Inl Inl o s1) s1') (case_sum (F2map id Inl Inl o s2)
s2') Inl Inl
  apply (tactic << rtac @context (thm mor_UNIV) RS iffD2) 1 >>)
  apply (rule conjI)
  apply (rule sym)
  apply (rule case_sum_o_inj(1))
  apply (rule sym)
  apply (rule case_sum_o_inj(1))
done

```

2.4 Bisimulations

definition bis where

```

  bis B1 B2 s1 s2 B1' B2' s1' s2' R1 R2 =
  ((R1 ⊆ B1 × B1' ∧ R2 ⊆ B2 × B2') ∧
  ((∀ b1 b1'. (b1, b1') ∈ R1 →
  (∃ z ∈ F1in UNIV R1 R2.
  F1map id fst fst z = s1 b1 ∧ F1map id snd snd z = s1' b1')) ∧
  (∀ b2 b2'. (b2, b2') ∈ R2 →
  (∃ z ∈ F2in UNIV R1 R2.
  F2map id fst fst z = s2 b2 ∧ F2map id snd snd z = s2' b2')))))

```

lemma bis_cong: $\llbracket \text{bis } B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ R1\ R2; R1' = R1; R2' = R2 \rrbracket \implies$

```

  bis B1 B2 s1 s2 B1' B2' s1' s2' R1' R2'
  apply (tactic << hyp_subst_tac @context 1 >>)
  apply assumption

```

done

lemma *bis_Frel*:

$bis\ B1\ B2\ s1\ s2\ B1'\ B2'\ s1'\ s2'\ R1\ R2 \longleftrightarrow$
 $(R1 \subseteq B1 \times B1' \wedge R2 \subseteq B2 \times B2') \wedge$
 $((\forall b1\ b1'. (b1, b1') \in R1 \longrightarrow F1rel (=) (in_rel\ R1)\ (in_rel\ R2)\ (s1\ b1)\ (s1'\ b1')) \wedge$
 $(\forall b2\ b2'. (b2, b2') \in R2 \longrightarrow F2rel (=) (in_rel\ R1)\ (in_rel\ R2)\ (s2\ b2)\ (s2'\ b2'))))$

apply (rule trans[OF bis_def])

apply (rule iffI)

apply (erule conjE)

apply (erule conjI)

apply (rule conjI)

apply (rule allI)

apply (rule allI)

apply (rule impI)

apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)

apply (erule allE)+

apply (erule impE)

apply assumption

apply (erule bexE)

apply (erule conjE CollectE)+

apply (rule iffD2[OF F1.in_rel])

apply (rule exI)

apply (rule conjI[rotated])

apply (rule conjI)

apply (rule trans)

apply (rule trans)

apply (rule F1.map_comp)

apply (rule F1.map_cong0)

apply (rule fst_diag_id)

apply (rule fun_cong[OF o_id])

apply (rule fun_cong[OF o_id])

apply assumption

apply (rule trans)

apply (rule trans)

apply (rule F1.map_comp)

apply (rule F1.map_cong0)

apply (rule snd_diag_id)

apply (rule fun_cong[OF o_id])

apply (rule fun_cong[OF o_id])

apply assumption

apply (rule CollectI)

apply (rule conjI)

apply (rule ord_eq_le_trans)

apply (rule F1.set_map(1))

apply (rule subset_trans)

apply (erule image_mono)

apply (rule image_subsetI)

apply (rule CollectI)

apply (rule case_prodI)

apply (rule refl)

apply (rule conjI)

apply (rule ord_eq_le_trans)

apply (rule trans)

apply (rule F1.set_map(2))

apply (rule trans[OF fun_cong[OF image_id] id_apply])

apply (erule Collect_case_prod_in_rel_leI)

apply (rule ord_eq_le_trans)

apply (rule trans)

apply (rule F1.set_map(3))

```

apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply (erule Collect_case_prod_in_rel_leI)

apply (rule allI)
apply (rule allI)
apply (rule impI)
apply (tactic ⟨ dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 ⟩)
apply (erule allE)+
apply (erule impE)
  apply assumption
apply (erule bexE)
apply (erule conjE CollectE)+
apply (rule iffD2[OF F2.in_rel])
apply (rule exI)
apply (rule conjI[rotated])
apply (rule conjI)
  apply (rule trans)
    apply (rule trans)
      apply (rule F2.map_comp)
      apply (rule F2.map_cong0)
        apply (rule fst_diag_id)
        apply (rule fun_cong[OF o_id])
        apply (rule fun_cong[OF o_id])
      apply assumption

apply (rule trans)
apply (rule trans)
  apply (rule F2.map_comp)
apply (rule F2.map_cong0)
  apply (rule snd_diag_id)
  apply (rule fun_cong[OF o_id])
apply (rule fun_cong[OF o_id])
apply assumption

apply (rule CollectI)
apply (rule conjI)
apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(1))
apply (rule subset_trans)
  apply (erule image_mono)
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule case_prodI)
apply (rule refl)
apply (rule conjI)
apply (rule ord_eq_le_trans)
  apply (rule trans)
    apply (rule F2.set_map(2))
    apply (rule trans[OF fun_cong[OF image_id] id_apply])
  apply (erule Collect_case_prod_in_rel_leI)
apply (rule ord_eq_le_trans)
apply (rule trans)
  apply (rule F2.set_map(3))
  apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply (erule Collect_case_prod_in_rel_leI)

apply (erule conjE)
apply (erule conjI)

apply (rule conjI)
apply (tactic ⟨ dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 ⟩)
apply (rule allI)
apply (rule allI)

```

```

apply (rule impI)
apply (erule allE)
apply (erule allE)
apply (erule impE)
  apply assumption
apply (drule iffD1[OF F1.in_rel])
apply (erule exE conjE CollectE Collect_case_prod_in_rel_leE)+

```

```

apply (rule beXI)
apply (rule conjI)
apply (rule trans)
  apply (rule F1.map_comp)
apply (tactic << stac @{context} @{thm id_o} 1 >>)
apply (tactic << stac @{context} @{thm o_id} 1 >>)
apply (tactic << stac @{context} @{thm o_id} 1 >>)
apply assumption

```

```

apply (rule trans)
  apply (rule F1.map_comp)
apply (tactic << stac @{context} @{thm id_o} 1 >>)
apply (tactic << stac @{context} @{thm o_id} 1 >>)
apply (tactic << stac @{context} @{thm o_id} 1 >>)
apply (rule trans)
  apply (rule F1.map_cong0)
    apply (rule Collect_case_prodD)
    apply (erule set_mp)
    apply assumption
    apply (rule refl)
  apply (rule refl)
apply assumption

```

```

apply (rule CollectI)
apply (rule conjI)
apply (rule subset_UNIV)

```

```

apply (rule conjI)
apply (rule ord_eq_le_trans)
  apply (rule trans)
    apply (rule F1.set_map(2))
  apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply assumption

```

```

apply (rule ord_eq_le_trans)
apply (rule trans)
  apply (rule F1.set_map(3))
apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply assumption

```

```

apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (rule allI)
apply (rule allI)
apply (rule impI)
apply (erule allE)
apply (erule allE)
apply (erule impE)
  apply assumption
apply (drule iffD1[OF F2.in_rel])
apply (erule exE conjE CollectE Collect_case_prod_in_rel_leE)+

```

```

apply (rule beXI)
apply (rule conjI)

```

```

apply (rule trans)
  apply (rule F2.map_comp)
apply (tactic << stac @ {context} @ {thm id_o} 1 >>)
apply (tactic << stac @ {context} @ {thm o_id} 1 >>)
apply (tactic << stac @ {context} @ {thm o_id} 1 >>)
apply assumption

apply (rule trans)
  apply (rule F2.map_comp)
apply (tactic << stac @ {context} @ {thm id_o} 1 >>)
apply (tactic << stac @ {context} @ {thm o_id} 1 >>)
apply (tactic << stac @ {context} @ {thm o_id} 1 >>)
apply (rule trans)
  apply (rule F2.map_cong0)
    apply (rule Collect_case_prodD)
    apply (erule set_mp)
    apply assumption
  apply (rule refl)
apply (rule refl)
apply assumption

apply (rule CollectI)
apply (rule conjI)
apply (rule subset_UNIV)

apply (rule conjI)
apply (rule ord_eq_le_trans)
  apply (rule trans)
    apply (rule F2.set_map(2))
  apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply assumption

apply (rule ord_eq_le_trans)
apply (rule trans)
  apply (rule F2.set_map(3))
apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply assumption
done

lemma bis_converse:
  bis B1 B2 s1 s2 B1' B2' s1' s2' R1 R2 ==>
  bis B1' B2' s1' s2' B1 B2 s1 s2 (R1^-1) (R2^-1)
apply (tactic << rtac @ {context} (@ {thm bis_Frel} RS iffD2) 1 >>)
apply (tactic << dtac @ {context} (@ {thm bis_Frel} RS iffD1) 1 >>)
apply (erule conjE)+
apply (rule conjI)

apply (rule conjI)
  apply (rule iffD1[OF converse_subset_swap])
  apply (erule subset_trans)
  apply (rule equalityD2)
  apply (rule converse_Times)

apply (rule iffD1[OF converse_subset_swap])
apply (erule subset_trans)
apply (rule equalityD2)
apply (rule converse_Times)

apply (rule conjI)
apply (rule allI)
apply (rule allI)
apply (rule impI)
apply (rule predicate2D[OF eq_refl[OF F1rel_cong]])

```

```

apply (rule conversep_eq)
apply (rule conversep_in_rel)
apply (rule conversep_in_rel)
apply (rule predicate2D[OF eq_refl[OF sym[OF F1.rel_conversep]]])
apply (erule allE)+
apply (rule conversepI)
apply (erule mp)
apply (erule converseD)

```

```

apply (rule allI)
apply (rule allI)
apply (rule impI)
apply (rule predicate2D[OF eq_refl[OF F2rel_cong]])
  apply (rule conversep_eq)
  apply (rule conversep_in_rel)
  apply (rule conversep_in_rel)
apply (rule predicate2D[OF eq_refl[OF sym[OF F2.rel_conversep]]])
apply (erule allE)+
apply (rule conversepI)
apply (erule mp)
apply (erule converseD)
done

```

lemma bis_Comp:

```

[[bis B1 B2 s1 s2 B1' B2' s1' s2' P1 P2;
  bis B1' B2' s1' s2' B1'' B2'' s1'' s2'' Q1 Q2]] ==>
  bis B1 B2 s1 s2 B1'' B2'' s1'' s2'' (P1 O Q1) (P2 O Q2)
apply (tactic << rtac @{context} (@{thm bis_Frel[THEN iffD2]} 1) >>)
apply (tactic << dtac @{context} (@{thm bis_Frel[THEN iffD1]} 1) >>)+
apply (erule conjE)+
apply (rule conjI)
apply (rule conjI)
  apply (erule relcomp_subset_Sigma)
  apply assumption
apply (erule relcomp_subset_Sigma)
apply assumption

```

```

apply (rule conjI)
apply (rule allI)+
apply (rule impI)
apply (rule predicate2D[OF eq_refl[OF F1rel_cong]])
  apply (rule eq_OO)
  apply (rule relcompp_in_rel)
  apply (rule relcompp_in_rel)
apply (rule predicate2D[OF eq_refl[OF sym[OF F1.rel_compp]]])
apply (erule relcompE)
apply (tactic << dtac @{context} (@{thm prod.inject[THEN iffD1]} 1) >>)
apply (erule conjE)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (erule allE)+
apply (rule relcomppI)
  apply (erule mp)
  apply assumption
apply (erule mp)
apply assumption

```

```

apply (rule allI)+
apply (rule impI)
apply (rule predicate2D[OF eq_refl[OF F2rel_cong]])
  apply (rule eq_OO)
  apply (rule relcompp_in_rel)
  apply (rule relcompp_in_rel)
apply (rule predicate2D[OF eq_refl[OF sym[OF F2.rel_compp]]])

```

```

apply (erule relcompE)
apply (tactic << dtac @{{context}} (@{thm prod.inject[THEN iffD1]}) 1 >>)
apply (erule conjE)
apply (tactic << hyp_subst_tac @{{context}} 1 >>)
apply (erule allE)+
apply (rule relcomppI)
  apply (erule mp)
  apply assumption
apply (erule mp)
apply assumption
done

```

lemma bis_Gr: $\llbracket \text{coalg } B1 \ B2 \ s1 \ s2; \text{mor } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ f1 \ f2 \rrbracket \implies$
 $\text{bis } B1 \ B2 \ s1 \ s2 \ B1' \ B2' \ s1' \ s2' \ (\text{BNF_Def.Gr } B1 \ f1) \ (\text{BNF_Def.Gr } B2 \ f2)$

```

unfolding bis_Frel eq_alt in_rel_Gr F1.rel_Grp F2.rel_Grp
apply (rule conjI)
apply (rule conjI)
  apply (rule iffD2[OF Gr_incl])
  apply (erule mor_image1)
apply (rule iffD2[OF Gr_incl])
apply (erule mor_image2)

```

```

apply (rule conjI)
apply (rule allI)
apply (rule allI)
apply (rule impI)
apply (rule GrpI)
  apply (erule trans[OF morE1])
  apply (erule GrD1)
  apply (erule arg_cong[OF GrD2])
apply (erule coalg_F1in)
apply (erule GrD1)

```

```

apply (rule allI)
apply (rule allI)
apply (rule impI)
apply (rule GrpI)
  apply (erule trans[OF morE2])
  apply (erule GrD1)
  apply (erule arg_cong[OF GrD2])
apply (erule coalg_F2in)
apply (erule GrD1)
done

```

lemmas bis_image2 = bis_cong[OF bis_Comp[OF bis_converse[OF bis_Gr] bis_Gr] image2_Gr image2_Gr]
lemmas bis_diag = bis_cong[OF bis_Gr[OF _ mor_id] Id_on_Gr Id_on_Gr]

lemma bis_Union: $\forall i \in I. \text{bis } B1 \ B2 \ s1 \ s2 \ B1 \ B2 \ s1 \ s2 \ (R1i \ i) \ (R2i \ i) \implies$
 $\text{bis } B1 \ B2 \ s1 \ s2 \ B1 \ B2 \ s1 \ s2 \ (\bigcup_{i \in I} R1i \ i) \ (\bigcup_{i \in I} R2i \ i)$

```

unfolding bis_def
apply (rule conjI)
apply (rule conjI)
  apply (rule UN_least)
  apply (erule bspec)
  apply assumption
  apply (erule conjunct1)
apply (tactic << etac @{{context}} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (rule UN_least)
apply (erule bspec)
  apply assumption
  apply (erule conjunct1)
apply (tactic << etac @{{context}} (BNF_Util.mk_conjunctN 2 2) 1 >>)

```

```

apply (rule conjI)
apply (rule allI)+
apply (rule impI)
apply (erule UN_E)
apply (drule bspec)
apply assumption
apply (drule conjunct2)
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule allE)+
apply (drule mp)
apply assumption
apply (erule bexE)
apply (rule bexE)
apply assumption
apply (rule F1in_mono23')
apply (erule SUP_upper2[OF _ subset_refl])
apply (erule SUP_upper2[OF _ subset_refl])
apply assumption

```

```

apply (rule allI)+
apply (rule impI)
apply (erule UN_E)
apply (drule bspec)
apply assumption
apply (drule conjunct2)
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule allE)+
apply (drule mp)
apply assumption
apply (erule bexE)
apply (rule bexE)
apply assumption
apply (rule F2in_mono23')
apply (erule SUP_upper2[OF _ subset_refl])
apply (erule SUP_upper2[OF _ subset_refl])
apply assumption
done

```

abbreviation $sbis\ B1\ B2\ s1\ s2\ R1\ R2 \equiv bis\ B1\ B2\ s1\ s2\ B1\ B2\ s1\ s2\ R1\ R2$

definition $lsbis1$ **where** $lsbis1\ B1\ B2\ s1\ s2 =$
 $(\bigcup R \in \{(R1, R2) \mid R1\ R2 . sbis\ B1\ B2\ s1\ s2\ R1\ R2\}. fst\ R)$

definition $lsbis2$ **where** $lsbis2\ B1\ B2\ s1\ s2 =$
 $(\bigcup R \in \{(R1, R2) \mid R1\ R2 . sbis\ B1\ B2\ s1\ s2\ R1\ R2\}. snd\ R)$

lemma $sbis_lsbis$:
 $sbis\ B1\ B2\ s1\ s2\ (lsbis1\ B1\ B2\ s1\ s2)\ (lsbis2\ B1\ B2\ s1\ s2)$
apply (tactic << rtac @{context} (Thm.permute_prem0 1 @{thm bis_cong}) 1 >>)
apply (rule lsbis1_def)
apply (rule lsbis2_def)
apply (rule bis_Union)
apply (rule ballI)
apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (erule bis_cong)
apply (rule fst_conv)
apply (rule snd_conv)
done

lemmas $lsbis1_incl = conjunct1[OF conjunct1[OF iffD1[OF bis_def]], OF sbis_lsbis]$


```

lemmas lsbis2_incl = conjunct2[OF conjunct1[OF iffD1[OF bis_def]], OF sbis_lsbis]
lemmas lsbisE1 =
  mp[OF spec[OF spec[OF conjunct1[OF conjunct2[OF iffD1[OF bis_def]], OF sbis_lsbis]]]
lemmas lsbisE2 =
  mp[OF spec[OF spec[OF conjunct2[OF conjunct2[OF iffD1[OF bis_def]], OF sbis_lsbis]]]

lemma incl_lsbis1: sbis B1 B2 s1 s2 R1 R2  $\implies$   $R1 \subseteq$  lsbis1 B1 B2 s1 s2
apply (rule xt1(3))
apply (rule lsbis1_def)
apply (rule SUP_upper2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply assumption
apply (rule equalityD2)
apply (rule fst_conv)
done

lemma incl_lsbis2: sbis B1 B2 s1 s2 R1 R2  $\implies$   $R2 \subseteq$  lsbis2 B1 B2 s1 s2
apply (rule xt1(3))
apply (rule lsbis2_def)
apply (rule SUP_upper2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply assumption
apply (rule equalityD2)
apply (rule snd_conv)
done

lemma equiv_lsbis1: coalg B1 B2 s1 s2  $\implies$  equiv B1 (lsbis1 B1 B2 s1 s2)
apply (rule iffD2[OF equiv_def])

apply (rule conjI)
apply (rule iffD2[OF refl_on_def])
apply (rule conjI)
apply (rule lsbis1_incl)
apply (rule ballI)
apply (rule set_mp)
apply (rule incl_lsbis1)
apply (rule bis_diag)
apply assumption
apply (erule Id_onI)

apply (rule conjI)
apply (rule iffD2[OF sym_def])
apply (rule allI)+
apply (rule impI)
apply (rule set_mp)
apply (rule incl_lsbis1)
apply (rule bis_converse)
apply (rule sbis_lsbis)
apply (erule converseI)

apply (rule iffD2[OF trans_def])
apply (rule allI)+
apply (rule impI)+
apply (rule set_mp)
apply (rule incl_lsbis1)
apply (rule bis_Comp)
apply (rule sbis_lsbis)

```

```

  apply (rule sbis_lsbis)
  apply (erule relcompI)
  apply assumption
done

```

lemma *equiv_lsbis2*: $\text{coalg } B1 \ B2 \ s1 \ s2 \implies \text{equiv } B2 \ (\text{lsbis2 } B1 \ B2 \ s1 \ s2)$
unfolding *equiv_def refl_on_def sym_def trans_def*
apply (rule *conjI*)

```

  apply (rule conjI)
  apply (rule lsbis2_incl)
  apply (rule ballI)
  apply (rule set_mp)
  apply (rule incl_lsbis2)
  apply (rule bis_diag)
  apply assumption
  apply (erule Id_onI)

```

```

  apply (rule conjI)

```

```

  apply (rule allI)+
  apply (rule impI)
  apply (rule set_mp)
  apply (rule incl_lsbis2)
  apply (rule bis_converse)
  apply (rule sbis_lsbis)
  apply (erule converseI)

```

```

  apply (rule allI)+
  apply (rule impI)+
  apply (rule set_mp)
  apply (rule incl_lsbis2)
  apply (rule bis_Comp)
  apply (rule sbis_lsbis)
  apply (rule sbis_lsbis)
  apply (erule relcompI)
  apply assumption
done

```

2.5 The Tree Coalgebra

```

typedef bd_type_F = UNIV :: (bd_type_F1 + bd_type_F2) set
  apply (rule exI) apply (rule UNIV_I)
  done

```

```

type-synonym 'a carrier = ((bd_type_F + bd_type_F) list set ×
  ((bd_type_F + bd_type_F) list ⇒ ('a, bd_type_F, bd_type_F) F1 + ('a, bd_type_F, bd_type_F) F2))

```

```

abbreviation bd_F ≡ dir_image (bd_F1 +c bd_F2) Abs_bd_type_F
lemmas bd_F = dir_image[OF Abs_bd_type_F_inject[OF UNIV_I UNIV_I] Card_order_csum]
lemmas bd_F_Cinfinite = Cinfinite_cong[OF bd_F_Cinfinite_csum1[OF F1.bd_Cinfinite]]
lemmas bd_F_Card_order = Card_order_ordIso[OF Card_order_csum_ordIso_symmetric[OF bd_F]]
lemma bd_F_card_order: card_order bd_F
  apply (rule card_order_dir_image)
  apply (rule bijI)
  apply (rule Abs_bd_type_F_inject[OF UNIV_I UNIV_I])
  apply (rule Abs_bd_type_F_cases)
  apply (erule exI)
  apply (rule card_order_csum)
  apply (rule F1.bd_card_order)
  apply (rule F2.bd_card_order)
  done

```

```

lemmas F1set1_bd' = ordLeq_transitive[OF F1.set_bd(1) ordLeq_ordIso_trans[OF ordLeq_csum1[OF F1.bd_Card_order]]

```

bd_F]]
lemmas *F1set2_bd'* = *ordLeq_transitive*[*OF F1.set_bd*(2) *ordLeq_ordIso_trans*[*OF ordLeq_csum1*[*OF F1.bd_Card_order*]
bd_F]]
lemmas *F1set3_bd'* = *ordLeq_transitive*[*OF F1.set_bd*(3) *ordLeq_ordIso_trans*[*OF ordLeq_csum1*[*OF F1.bd_Card_order*]
bd_F]]

lemmas *F2set1_bd'* = *ordLeq_transitive*[*OF F2.set_bd*(1) *ordLeq_ordIso_trans*[*OF ordLeq_csum2*[*OF F2.bd_Card_order*]
bd_F]]
lemmas *F2set2_bd'* = *ordLeq_transitive*[*OF F2.set_bd*(2) *ordLeq_ordIso_trans*[*OF ordLeq_csum2*[*OF F2.bd_Card_order*]
bd_F]]
lemmas *F2set3_bd'* = *ordLeq_transitive*[*OF F2.set_bd*(3) *ordLeq_ordIso_trans*[*OF ordLeq_csum2*[*OF F2.bd_Card_order*]
bd_F]]

abbreviation *Succ1 Kl kl* \equiv {*k1*. *Inl k1* \in *BNF_Greatest_Fixpoint.Succ Kl kl*}
abbreviation *Succ2 Kl kl* \equiv {*k2*. *Inr k2* \in *BNF_Greatest_Fixpoint.Succ Kl kl*}

definition *isNode1* **where**
isNode1 Kl lab kl = ($\exists x1$. *lab kl* = *Inl x1* \wedge *F1set2 x1* = *Succ1 Kl kl* \wedge *F1set3 x1* = *Succ2 Kl kl*)

definition *isNode2* **where**
isNode2 Kl lab kl = ($\exists x2$. *lab kl* = *Inr x2* \wedge *F2set2 x2* = *Succ1 Kl kl* \wedge *F2set3 x2* = *Succ2 Kl kl*)

abbreviation *isTree* **where**
isTree Kl lab \equiv ($[] \in Kl \wedge$
 $(\forall kl \in Kl. (\forall k1 \in Succ1 Kl kl. isNode1 Kl lab (kl @ [Inl k1])) \wedge$
 $(\forall k2 \in Succ2 Kl kl. isNode2 Kl lab (kl @ [Inr k2])))$)

definition *carT1* **where**
carT1 = {(*Kl* :: (*bd_type_F* + *bd_type_F*) *list set*, *lab*) | *Kl lab. isTree Kl lab* \wedge *isNode1 Kl lab []*}

definition *carT2* **where**
carT2 = {(*Kl* :: (*bd_type_F* + *bd_type_F*) *list set*, *lab*) | *Kl lab. isTree Kl lab* \wedge *isNode2 Kl lab []*}

definition *strT1* **where**
strT1 = (*case_prod* (%*Kl lab. case_sum* (*F1map id*
 $(\lambda k1. (BNF_Greatest_Fixpoint.Shift Kl (Inl k1), BNF_Greatest_Fixpoint.shift lab (Inl k1)))$
 $(\lambda k2. (BNF_Greatest_Fixpoint.Shift Kl (Inr k2), BNF_Greatest_Fixpoint.shift lab (Inr k2)))$) *undefined* (*lab*
 $[]$)))

definition *strT2* **where**
strT2 = (*case_prod* (%*Kl lab. case_sum* *undefined* (*F2map id*
 $(\lambda k1. (BNF_Greatest_Fixpoint.Shift Kl (Inl k1), BNF_Greatest_Fixpoint.shift lab (Inl k1)))$
 $(\lambda k2. (BNF_Greatest_Fixpoint.Shift Kl (Inr k2), BNF_Greatest_Fixpoint.shift lab (Inr k2)))$) (*lab []*)))

lemma *coalg_T*: *coalg carT1 carT2 strT1 strT2*
unfolding *coalg_def carT1_def carT2_def isNode1_def isNode2_def*
apply (*rule conjI*)
apply (*rule ballI*)
apply (*erule CollectE exE conjE*)
apply (*tactic* $\langle\langle hyp_subst_tac @\{context\} 1 \rangle\rangle$)
apply (*rule ssubst_mem*[*OF trans*[*OF trans*[*OF fun_cong*[*OF strT1_def*] *prod.case*]]])
apply (*erule trans*[*OF arg_cong*])
apply (*rule sum.case*(1))
apply (*rule CollectI*)
apply (*rule conjI*)
apply (*rule subset_UNIV*)
apply (*rule conjI*)
apply (*rule ord_eq_le_trans*[*OF F1.set_map*(2)])
apply (*rule image_subsetI*)
apply (*rule CollectI*)
apply (*rule exI*)
apply (*rule conjI*)
apply (*rule refl*)

```

apply (rule conjI)
apply (rule conjI)
  apply (erule empty_Shift)
  apply (drule set_rev_mp)
    apply (erule equalityD1)
  apply (erule CollectD)
apply (rule ballI)
apply (rule conjI)
  apply (rule ballI)
  apply (erule CollectE)
  apply (drule ShiftD)
  apply (drule bspec)
  apply (erule thin_rl)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule bspec)
  apply (rule CollectI)
  apply (erule set_mp[OF equalityD1[OF Succ_Shift]])
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
  apply (rule trans[OF fun_cong[OF shift_def]])
  apply (rule trans[OF arg_cong[OF sym[OF append_Cons]]])
  apply assumption
apply (rule conjI)
  apply (erule trans)
  apply (rule Collect_cong)
  apply (rule eqset_imp_iff)
  apply (rule sym)
  apply (rule trans)
  apply (rule Succ_Shift)
  apply (rule arg_cong[OF sym[OF append_Cons]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])

```

```

apply (rule ballI)
apply (erule CollectE)
apply (drule ShiftD)
apply (drule bspec)
apply (erule thin_rl)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule bspec)
  apply (rule CollectI)
  apply (erule set_mp[OF equalityD1[OF Succ_Shift]])
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
  apply (rule trans[OF fun_cong[OF shift_def]])
  apply (rule trans[OF arg_cong[OF sym[OF append_Cons]]])
  apply assumption
apply (rule conjI)
apply (erule trans)
  apply (rule Collect_cong)
  apply (rule eqset_imp_iff)
  apply (rule sym)
apply (rule trans)

```

```

    apply (rule Succ_Shift)
    apply (rule arg_cong[OF sym[OF append_Cons]])
  apply (erule trans)
  apply (rule Collect_cong)
  apply (rule eqset_imp_iff)
  apply (rule sym)
  apply (rule trans)
  apply (rule Succ_Shift)
  apply (rule arg_cong[OF sym[OF append_Cons]])

  apply (drule bspec)
  apply assumption
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
  apply (drule bspec)
  apply (erule set_mp[OF equalityD1])
  apply assumption
  apply (erule exE conjE)+
  apply (rule exI)
  apply (rule conjI)
  apply (rule trans[OF fun_cong[OF shift_def]])
  apply (erule trans[OF arg_cong[OF sym[OF append_Nil]]])
  apply (rule conjI)
  apply (erule trans)
  apply (rule Collect_cong)
  apply (rule eqset_imp_iff)
  apply (rule sym)
  apply (rule trans)
  apply (rule Succ_Shift)
  apply (rule arg_cong[OF sym[OF append_Nil]])
  apply (erule trans)
  apply (rule Collect_cong)
  apply (rule eqset_imp_iff)
  apply (rule sym)
  apply (rule trans)
  apply (rule Succ_Shift)
  apply (rule arg_cong[OF sym[OF append_Nil]])

  apply (rule ord_eq_le_trans[OF F1.set_map(3)])
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule exI)+
  apply (rule conjI)
  apply (rule refl)
  apply (rule conjI)
  apply (rule conjI)
  apply (erule empty_Shift)
  apply (drule set_rev_mp)
  apply (erule equalityD1)
  apply (erule CollectD)
  apply (rule ballI)
  apply (rule conjI)
  apply (rule ballI)
  apply (erule CollectE)
  apply (drule ShiftD)
  apply (drule bspec)
  apply (erule thin_rl)
  apply assumption
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
  apply (drule bspec)
  apply (rule CollectI)
  apply (erule set_mp[OF equalityD1[OF Succ_Shift]])
  apply (erule exE conjE)+
  apply (rule exI)

```

```

apply (rule conjI)
apply (rule trans[OF fun_cong[OF shift_def]])
apply (rule trans[OF arg_cong[OF sym[OF append_Cons]]])
apply assumption
apply (rule conjI)
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])

apply (rule ballI)
apply (erule CollectE)
apply (drule ShiftD)
apply (drule bspec)
  apply (erule thin_rl)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule bspec)
  apply (rule CollectI)
  apply (erule set_mp[OF equalityD1[OF Succ_Shift]])
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
  apply (rule trans[OF fun_cong[OF shift_def]])
  apply (rule trans[OF arg_cong[OF sym[OF append_Cons]]])
apply assumption
apply (rule conjI)
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])

apply (drule bspec)
apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule bspec)
  apply (erule set_mp[OF equalityD1])
  apply assumption
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
apply (rule trans[OF fun_cong[OF shift_def]])

```

```

apply (erule trans[OF arg_cong[OF sym[OF append_Nil]])
apply (rule conjI)
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Nil]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Nil]])

```

```

apply (rule ballI)
apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (rule ssubst_mem[OF trans[OF fun_cong[OF strT2_def] prod.case]])
apply (rule ssubst_mem)
apply (rule trans)
  apply (erule arg_cong)
  apply (rule sum.case(2))
apply (rule CollectI)
apply (rule conjI)
apply (rule subset_UNIV)
apply (rule conjI)
apply (rule ord_eq_le_trans[OF F2.set_map(2)])
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply (rule conjI)
apply (rule conjI)
  apply (erule empty_Shift)
  apply (drule set_rev_mp)
  apply (erule equalityD1)
  apply (erule CollectD)
apply (rule ballI)
apply (rule conjI)
apply (rule ballI)
apply (erule CollectE)
apply (drule ShiftD)
apply (drule bspec)
  apply (erule thin_rl)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule bspec)
  apply (rule CollectI)
  apply (erule set_mp[OF equalityD1[OF Succ_Shift]])
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
apply (rule trans[OF fun_cong[OF shift_def]])
  apply (rule trans[OF arg_cong[OF sym[OF append_Cons]])
  apply assumption
apply (rule conjI)
apply (erule trans)

```

```

apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])

```

```

apply (rule ballI)
apply (erule CollectE)
apply (drule ShiftD)
apply (drule bspec)
  apply (erule thin_rl)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule bspec)
  apply (rule CollectI)
  apply (erule set_mp[OF equalityD1[OF Succ_Shift]])
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
  apply (rule trans[OF fun_cong[OF shift_def]])
  apply (rule trans[OF arg_cong[OF sym[OF append_Cons]]])
apply assumption
apply (rule conjI)
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])

```

```

apply (drule bspec)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule bspec)
  apply (erule set_mp[OF equalityD1])
apply assumption
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
  apply (rule trans[OF fun_cong[OF shift_def]])
  apply (erule trans[OF arg_cong[OF sym[OF append_Nil]]])
apply (rule conjI)
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)

```



```

apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Nil]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Nil]])

apply (rule ord_eq_le_trans[OF F2.set_map(3)])
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (rule conjI)
apply (rule conjI)
  apply (erule empty_Shift)
  apply (drule set_rev_mp)
  apply (erule equalityD1)
  apply (erule CollectD)
apply (rule ballI)
apply (rule conjI)
apply (rule ballI)
apply (erule CollectE)
apply (drule ShiftD)
apply (drule bspec)
  apply (erule thin_rl)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>))
apply (drule bspec)
  apply (rule CollectI)
apply (erule set_mp[OF equalityD1[OF Succ_Shift]])
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
  apply (rule trans[OF fun_cong[OF shift_def]])
  apply (rule trans[OF arg_cong[OF sym[OF append_Cons]]])
apply assumption
apply (rule conjI)
  apply (erule trans)
  apply (rule Collect_cong)
  apply (rule eqset_imp_iff)
  apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
  apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])

apply (rule ballI)
apply (erule CollectE)
apply (drule ShiftD)
apply (drule bspec)
apply (erule thin_rl)

```

```

apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (drule bspec)
apply (rule CollectI)
apply (erule set_mp[OF equalityD1[OF Succ_Shift]])
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
apply (rule trans[OF fun_cong[OF shift_def]])
apply (rule trans[OF arg_cong[OF sym[OF append_Cons]]])
apply assumption
apply (rule conjI)
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Cons]])

```

```

apply (drule bspec)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (drule bspec)
apply (erule set_mp[OF equalityD1])
apply assumption
apply (erule exE conjE)+
apply (rule exI)
apply (rule conjI)
apply (rule trans[OF fun_cong[OF shift_def]])
apply (erule trans[OF arg_cong[OF sym[OF append_Nil]]])
apply (rule conjI)
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Nil]])
apply (erule trans)
apply (rule Collect_cong)
apply (rule eqset_imp_iff)
apply (rule sym)
apply (rule trans)
apply (rule Succ_Shift)
apply (rule arg_cong[OF sym[OF append_Nil]])
done

```

abbreviation *tobd_F12* **where** *tobd_F12 s1 x* \equiv *toCard* (*F1set2* (*s1 x*)) *bd_F*

abbreviation *tobd_F13* **where** *tobd_F13 s1 x* \equiv *toCard* (*F1set3* (*s1 x*)) *bd_F*

abbreviation *tobd_F22* **where** *tobd_F22 s2 x* \equiv *toCard* (*F2set2* (*s2 x*)) *bd_F*

abbreviation *tobd_F23* **where** *tobd_F23 s2 x* \equiv *toCard* (*F2set3* (*s2 x*)) *bd_F*

abbreviation *frombd_F12* **where** *frombd_F12 s1 x* \equiv *fromCard* (*F1set2* (*s1 x*)) *bd_F*

abbreviation *frombd_F13* **where** *frombd_F13 s1 x* \equiv *fromCard* (*F1set3* (*s1 x*)) *bd_F*

abbreviation *frombd_F22* **where** *frombd_F22 s2 x* \equiv *fromCard* (*F2set2* (*s2 x*)) *bd_F*

abbreviation *frombd_F23* **where** *frombd_F23 s2 x* \equiv *fromCard* (*F2set3* (*s2 x*)) *bd_F*

lemmas $\text{tobd_F12_inj} = \text{toCard_inj}[OF\ F1set2_bd'\ bd_F_Card_order]$
lemmas $\text{tobd_F13_inj} = \text{toCard_inj}[OF\ F1set3_bd'\ bd_F_Card_order]$
lemmas $\text{tobd_F22_inj} = \text{toCard_inj}[OF\ F2set2_bd'\ bd_F_Card_order]$
lemmas $\text{tobd_F23_inj} = \text{toCard_inj}[OF\ F2set3_bd'\ bd_F_Card_order]$
lemmas $\text{frombd_F12_tobd_F12} = \text{fromCard_toCard}[OF\ F1set2_bd'\ bd_F_Card_order]$
lemmas $\text{frombd_F13_tobd_F13} = \text{fromCard_toCard}[OF\ F1set3_bd'\ bd_F_Card_order]$
lemmas $\text{frombd_F22_tobd_F22} = \text{fromCard_toCard}[OF\ F2set2_bd'\ bd_F_Card_order]$
lemmas $\text{frombd_F23_tobd_F23} = \text{fromCard_toCard}[OF\ F2set3_bd'\ bd_F_Card_order]$

definition *Lev* where

$\text{Lev } s1\ s2 = \text{rec_nat } (\%a. \{\}, \%b. \{\})$
 $(\%n\ \text{rec.}$
 $(\%a1.$
 $\{ \text{Inl } (\text{tobd_F12 } s1\ a1\ b1) \# kl \mid b1\ kl. b1 \in F1set2\ (s1\ a1) \wedge kl \in \text{fst rec } b1 \} \cup$
 $\{ \text{Inr } (\text{tobd_F13 } s1\ a1\ b2) \# kl \mid b2\ kl. b2 \in F1set3\ (s1\ a1) \wedge kl \in \text{snd rec } b2 \},$
 $\%a2.$
 $\{ \text{Inl } (\text{tobd_F22 } s2\ a2\ b1) \# kl \mid b1\ kl. b1 \in F2set2\ (s2\ a2) \wedge kl \in \text{fst rec } b1 \} \cup$
 $\{ \text{Inr } (\text{tobd_F23 } s2\ a2\ b2) \# kl \mid b2\ kl. b2 \in F2set3\ (s2\ a2) \wedge kl \in \text{snd rec } b2 \})$

abbreviation *Lev1* where $\text{Lev1 } s1\ s2\ n \equiv \text{fst } (\text{Lev } s1\ s2\ n)$

abbreviation *Lev2* where $\text{Lev2 } s1\ s2\ n \equiv \text{snd } (\text{Lev } s1\ s2\ n)$

lemmas $\text{Lev1_0} = \text{fun_cong}[OF\ \text{fstI}[OF\ \text{rec_nat_0_imp}[OF\ \text{Lev_def}]]]$
lemmas $\text{Lev2_0} = \text{fun_cong}[OF\ \text{sndI}[OF\ \text{rec_nat_0_imp}[OF\ \text{Lev_def}]]]$
lemmas $\text{Lev1_Suc} = \text{fun_cong}[OF\ \text{fstI}[OF\ \text{rec_nat_Suc_imp}[OF\ \text{Lev_def}]]]$
lemmas $\text{Lev2_Suc} = \text{fun_cong}[OF\ \text{sndI}[OF\ \text{rec_nat_Suc_imp}[OF\ \text{Lev_def}]]]$

definition *rv* where

$\text{rv } s1\ s2 = \text{rec_list } (\%b1. \text{Inl } b1, \%b2. \text{Inr } b2)$
 $(\%k\ kl\ \text{rec.}$
 $(\%b1. \text{case_sum } (\%k1. \text{fst rec } (\text{frombd_F12 } s1\ b1\ k1)) (\%k2. \text{snd rec } (\text{frombd_F13 } s1\ b1\ k2))\ k,$
 $\%b2. \text{case_sum } (\%k1. \text{fst rec } (\text{frombd_F22 } s2\ b2\ k1)) (\%k2. \text{snd rec } (\text{frombd_F23 } s2\ b2\ k2))\ k)$

abbreviation *rv1* where $\text{rv1 } s1\ s2\ kl \equiv \text{fst } (\text{rv } s1\ s2\ kl)$

abbreviation *rv2* where $\text{rv2 } s1\ s2\ kl \equiv \text{snd } (\text{rv } s1\ s2\ kl)$

lemmas $\text{rv1_Nil} = \text{fun_cong}[OF\ \text{fstI}[OF\ \text{rec_list_Nil_imp}[OF\ \text{rv_def}]]]$
lemmas $\text{rv2_Nil} = \text{fun_cong}[OF\ \text{sndI}[OF\ \text{rec_list_Nil_imp}[OF\ \text{rv_def}]]]$
lemmas $\text{rv1_Cons} = \text{fun_cong}[OF\ \text{fstI}[OF\ \text{rec_list_Cons_imp}[OF\ \text{rv_def}]]]$
lemmas $\text{rv2_Cons} = \text{fun_cong}[OF\ \text{sndI}[OF\ \text{rec_list_Cons_imp}[OF\ \text{rv_def}]]]$

abbreviation *Lab1* $s1\ s2\ b1\ kl \equiv$

$(\text{case_sum } (\%k. \text{Inl } (F1map\ id\ (\text{tobd_F12 } s1\ k)\ (\text{tobd_F13 } s1\ k)\ (s1\ k)))$
 $(\%k. \text{Inr } (F2map\ id\ (\text{tobd_F22 } s2\ k)\ (\text{tobd_F23 } s2\ k)\ (s2\ k))) (\text{rv1 } s1\ s2\ kl\ b1))$

abbreviation *Lab2* $s1\ s2\ b2\ kl \equiv$

$(\text{case_sum } (\%k. \text{Inl } (F1map\ id\ (\text{tobd_F12 } s1\ k)\ (\text{tobd_F13 } s1\ k)\ (s1\ k)))$
 $(\%k. \text{Inr } (F2map\ id\ (\text{tobd_F22 } s2\ k)\ (\text{tobd_F23 } s2\ k)\ (s2\ k))) (\text{rv2 } s1\ s2\ kl\ b2))$

definition *beh1* $s1\ s2\ a = (\bigcup n. \text{Lev1 } s1\ s2\ n\ a, \text{Lab1 } s1\ s2\ a)$

definition *beh2* $s1\ s2\ a = (\bigcup n. \text{Lev2 } s1\ s2\ n\ a, \text{Lab2 } s1\ s2\ a)$

lemma *length_Lev*:

$\forall kl\ b1\ b2. (kl \in \text{Lev1 } s1\ s2\ n\ b1 \longrightarrow \text{length } kl = n) \wedge$
 $(kl \in \text{Lev2 } s1\ s2\ n\ b2 \longrightarrow \text{length } kl = n)$

apply (*rule nat_induct*)

apply (*rule allI*)⁺

apply (*rule conjI*)

```

apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev1_0]])
apply (erule singletonE)
apply (erule ssubst)
apply (rule list.size(3))

```

```

apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev2_0]])
apply (erule singletonE)
apply (erule ssubst)
apply (rule list.size(3))

```

```

apply (rule allI)+
apply (rule conjI)
apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev1_Suc]])
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule trans)
  apply (rule length_Cons)
apply (rule arg_cong[of _ _ Suc])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule trans)
  apply (rule length_Cons)
apply (rule arg_cong[of _ _ Suc])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev2_Suc]])
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule trans)
  apply (rule length_Cons)
apply (rule arg_cong[of _ _ Suc])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule trans)
  apply (rule length_Cons)
apply (rule arg_cong[of _ _ Suc])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply assumption
done

```

```

lemmas length_Lev1 = mp[OF conjunct1[OF spec[OF spec [OF spec[OF length_Lev]]]]]

```

lemmas *length_Lev2* = *mp[OF conjunct2[OF spec[OF spec [OF spec [OF length_Lev]]]]]*

lemma *length_Lev1'*: $kl \in Lev1\ s1\ s2\ n\ a \implies kl \in Lev1\ s1\ s2\ (length\ kl)\ a$
apply (*frule length_Lev1*)
apply (*erule ssubst*)
apply *assumption*
done

lemma *length_Lev2'*: $kl \in Lev2\ s1\ s2\ n\ a \implies kl \in Lev2\ s1\ s2\ (length\ kl)\ a$
apply (*frule length_Lev2*)
apply (*erule ssubst*)
apply *assumption*
done

lemma *rv_last*:

$\forall k\ b1\ b2.$

$((\exists b1'.\ rv1\ s1\ s2\ (kl\ @\ [Inl\ k])\ b1 = Inl\ b1') \wedge$
 $(\exists b1'.\ rv1\ s1\ s2\ (kl\ @\ [Inr\ k])\ b1 = Inr\ b1')) \wedge$
 $((\exists b2'.\ rv2\ s1\ s2\ (kl\ @\ [Inl\ k])\ b2 = Inl\ b2') \wedge$
 $(\exists b2'.\ rv2\ s1\ s2\ (kl\ @\ [Inr\ k])\ b2 = Inr\ b2'))$

apply (*rule list.induct[of _ kl]*)
apply (*rule allI*)+
apply (*rule conjI*)
apply (*rule conjI*)
apply (*rule exI*)
apply (*rule trans[OF arg_cong[OF append_Nil]]*)
apply (*rule trans[OF rv1_Cons]*)
apply (*rule trans[OF arg_cong[OF sum.case(1)]]*)
apply (*rule rv1_Nil*)
apply (*rule exI*)
apply (*rule trans[OF arg_cong[OF append_Nil]]*)
apply (*rule trans[OF rv1_Cons]*)
apply (*rule trans[OF arg_cong[OF sum.case(2)]]*)
apply (*rule rv2_Nil*)
apply (*rule conjI*)
apply (*rule exI*)
apply (*rule trans[OF arg_cong[OF append_Nil]]*)
apply (*rule trans[OF rv2_Cons]*)
apply (*rule trans[OF arg_cong[OF sum.case(1)]]*)
apply (*rule rv1_Nil*)
apply (*rule exI*)
apply (*rule trans[OF arg_cong[OF append_Nil]]*)
apply (*rule trans[OF rv2_Cons]*)
apply (*rule trans[OF arg_cong[OF sum.case(2)]]*)
apply (*rule rv2_Nil*)

apply (*rule allI*)+
apply (*rule sum.exhaust*)
apply (*rule conjI*)
apply (*erule allE*)+
apply (*tactic* $\ll\ dtac\ @\{context\}\ (BNF_Util.mk_conjunctN\ 2\ 1)\ 1\ \gg$)
apply (*rule conjI*)
apply (*tactic* $\ll\ dtac\ @\{context\}\ (BNF_Util.mk_conjunctN\ 2\ 1)\ 1\ \gg$)
apply (*erule exE*)
apply (*rule exI*)
apply (*rule trans[OF arg_cong[OF append_Cons]]*)
apply (*rule trans[OF rv1_Cons]*)
apply (*erule trans[OF arg_cong[OF sum.case_cong_weak]]*)
apply (*rule trans[OF arg_cong[OF sum.case(1)]]*)
apply *assumption*
apply (*tactic* $\ll\ dtac\ @\{context\}\ (BNF_Util.mk_conjunctN\ 2\ 2)\ 1\ \gg$)
apply (*erule exE*)
apply (*rule exI*)

```

apply (rule trans[OF arg_cong[OF append_Cons]])
apply (rule trans[OF rv1_Cons])
apply (erule trans[OF arg_cong[OF sum.case_cong_weak]])
apply (rule trans[OF arg_cong[OF sum.case(1)]])
apply assumption

apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (rule conjI)
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule exE)
apply (rule exI)
apply (rule trans[OF arg_cong[OF append_Cons]])
apply (rule trans[OF rv2_Cons])
apply (erule trans[OF arg_cong[OF sum.case_cong_weak]])
apply (rule trans[OF arg_cong[OF sum.case(1)]])
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule exE)
apply (rule exI)
apply (rule trans[OF arg_cong[OF append_Cons]])
apply (rule trans[OF rv2_Cons])
apply (erule trans[OF arg_cong[OF sum.case_cong_weak]])
apply (rule trans[OF arg_cong[OF sum.case(1)]])
apply assumption

apply (rule conjI)
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (rule conjI)
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule exE)
apply (rule exI)
apply (rule trans[OF arg_cong[OF append_Cons]])
apply (rule trans[OF rv1_Cons])
apply (erule trans[OF arg_cong[OF sum.case_cong_weak]])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule exE)
apply (rule exI)
apply (rule trans[OF arg_cong[OF append_Cons]])
apply (rule trans[OF rv1_Cons])
apply (erule trans[OF arg_cong[OF sum.case_cong_weak]])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply assumption

apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (rule conjI)
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule exE)
apply (rule exI)
apply (rule trans[OF arg_cong[OF append_Cons]])
apply (rule trans[OF rv2_Cons])
apply (erule trans[OF arg_cong[OF sum.case_cong_weak]])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule exE)
apply (rule exI)
apply (rule trans[OF arg_cong[OF append_Cons]])
apply (rule trans[OF rv2_Cons])

```

```

apply (erule trans[OF arg_cong[OF sum.case_cong_weak]])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply assumption
done

```

```

lemmas rv_last' = spec[OF spec[OF spec[OF rv_last]]]
lemmas rv1_Inl_last = conjunct1[OF conjunct1[OF rv_last']]
lemmas rv1_Inr_last = conjunct2[OF conjunct1[OF rv_last']]
lemmas rv2_Inl_last = conjunct1[OF conjunct2[OF rv_last']]
lemmas rv2_Inr_last = conjunct2[OF conjunct2[OF rv_last']]

```

lemma Fset_Lev:

```

 $\forall kl\ b1\ b2\ b1'\ b2'\ b1''\ b2''.$ 
(kl  $\in$  Lev1 s1 s2 n b1  $\longrightarrow$ 
  ((rv1 s1 s2 kl b1 = Inl b1'  $\longrightarrow$ 
    (b1''  $\in$  F1set2 (s1 b1')  $\longrightarrow$ 
      kl @ [Inl (tobd_F12 s1 b1' b1'')]  $\in$  Lev1 s1 s2 (Suc n) b1)  $\wedge$ 
      (b2''  $\in$  F1set3 (s1 b1')  $\longrightarrow$ 
        kl @ [Inr (tobd_F13 s1 b1' b2'')]  $\in$  Lev1 s1 s2 (Suc n) b1))  $\wedge$ 
    (rv1 s1 s2 kl b1 = Inr b2'  $\longrightarrow$ 
      (b1''  $\in$  F2set2 (s2 b2')  $\longrightarrow$ 
        kl @ [Inl (tobd_F22 s2 b2' b1'')]  $\in$  Lev1 s1 s2 (Suc n) b1)  $\wedge$ 
        (b2''  $\in$  F2set3 (s2 b2')  $\longrightarrow$ 
          kl @ [Inr (tobd_F23 s2 b2' b2'')]  $\in$  Lev1 s1 s2 (Suc n) b1))))  $\wedge$ 
  (kl  $\in$  Lev2 s1 s2 n b2  $\longrightarrow$ 
    ((rv2 s1 s2 kl b2 = Inl b1'  $\longrightarrow$ 
      (b1''  $\in$  F1set2 (s1 b1')  $\longrightarrow$ 
        kl @ [Inl (tobd_F12 s1 b1' b1'')]  $\in$  Lev2 s1 s2 (Suc n) b2)  $\wedge$ 
        (b2''  $\in$  F1set3 (s1 b1')  $\longrightarrow$ 
          kl @ [Inr (tobd_F13 s1 b1' b2'')]  $\in$  Lev2 s1 s2 (Suc n) b2))  $\wedge$ 
      (rv2 s1 s2 kl b2 = Inr b2'  $\longrightarrow$ 
        (b1''  $\in$  F2set2 (s2 b2')  $\longrightarrow$ 
          kl @ [Inl (tobd_F22 s2 b2' b1'')]  $\in$  Lev2 s1 s2 (Suc n) b2)  $\wedge$ 
          (b2''  $\in$  F2set3 (s2 b2')  $\longrightarrow$ 
            kl @ [Inr (tobd_F23 s2 b2' b2'')]  $\in$  Lev2 s1 s2 (Suc n) b2))))
apply (rule nat_induct[of _ n])

```

```

apply (rule allI)+
apply (rule conjI)
apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev1_0]])
apply (erule singletonE)
apply (tactic « hyp_subst_tac @ {context} 1 »)
apply (rule conjI)
apply (rule impI)
apply (drule trans[OF sym])
apply (rule rv1_Nil)
apply (drule Inl_inject)
apply (tactic « hyp_subst_tac @ {context} 1 »)
apply (rule conjI)
apply (rule impI)
apply (rule ssubst_mem[OF append_Nil])
apply (rule set_mp[OF equalityD2])
apply (rule Lev1_Suc)
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply (erule conjI)
apply (rule set_mp[OF equalityD2])
apply (rule Lev1_0)
apply (rule singletonI)

```

```

apply (rule impI)
apply (rule ssubst_mem[OF append_Nil])
apply (rule set_mp[OF equalityD2])
  apply (rule Lev1_Suc)
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_0)
apply (rule singletonI)

```

```

apply (rule impI)
apply (erule trans[OF sym])
  apply (rule rv1_Nil)
apply (erule notE[OF Inr_not_Inl])

```

```

apply (rule impI)
apply (erule set_rev_mp[OF _equalityD1])
  apply (rule Lev2_0)
apply (erule singletonE)
apply (tactic ⟨⟨ hyp_subst_tac @{context} 1 ⟩⟩)
apply (rule conjI)
  apply (rule impI)
  apply (erule trans[OF sym])
    apply (rule rv2_Nil)
apply (erule notE[OF Inl_not_Inr])

```

```

apply (rule impI)
apply (erule trans[OF sym])
  apply (rule rv2_Nil)
apply (erule Inr_inject)
apply (tactic ⟨⟨ hyp_subst_tac @{context} 1 ⟩⟩)
apply (tactic ⟨⟨ stac @{context} @{thm append_Nil} 1 ⟩⟩)+
apply (rule conjI)
  apply (rule impI)
  apply (rule set_mp[OF equalityD2])
    apply (rule Lev2_Suc)
  apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev1_0)
apply (rule singletonI)
apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_Suc)
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_0)
apply (rule singletonI)

```



```

apply (rule allI)+
apply (rule conjI)
apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev1_Suc]])
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (rule conjI)
apply (rule impI)
apply (rule set_mp[OF equalityD2])
apply (rule Lev1_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply (rule conjI)
apply assumption
apply (drule sym[OF trans[OF sym]])
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF sum.case(1)])
apply (erule arg_cong[OF frombd_F12_tobd_F12])
apply (erule allE)+
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply assumption

apply (rule impI)
apply (rule set_mp[OF equalityD2])
apply (rule Lev1_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply (rule conjI)
apply assumption
apply (drule sym[OF trans[OF sym]])
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF sum.case(1)])
apply (erule arg_cong[OF frombd_F12_tobd_F12])
apply (erule allE)+
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply assumption

apply (rule impI)

```

```

apply (rule conjI)
apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev1_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (rule conjI)
apply assumption
apply (drule sym[OF trans[OF sym]])
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF sum.case(1)])
apply (erule arg_cong[OF frombd_F12_tobd_F12])
apply (erule allE)+
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev1_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (rule conjI)
  apply assumption
apply (drule sym[OF trans[OF sym]])
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF sum.case(1)])
apply (erule arg_cong[OF frombd_F12_tobd_F12])
apply (erule allE)+
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (tactic << stac @ {context} @ {thm rv1_Cons} 1 >>)
apply (tactic << stac @ {context} @ {thm sum.case(2)} 1 >>)
apply (tactic << stac @ {context} @ {thm frombd_F13_tobd_F13} 1 >>)
apply assumption
apply (rule conjI)
apply (rule impI)
apply (rule conjI)
apply (rule impI)

```

```

apply (rule set_mp[OF equalityD2])
apply (rule Lev1_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev1_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (rule conjI)
apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev1_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev1_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic ⟨ dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 ⟩)
apply (drule mp)
  apply assumption
apply (tactic ⟨ dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 ⟩)
apply (drule mp)
  apply assumption
apply (tactic ⟨ dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 ⟩)
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (drule set_rev_mp[OF _ equalityD1])
  apply (rule Lev2_Suc)
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (tactic ⟨ hyp_subst_tac @ {context} 1 ⟩)
apply (tactic ⟨ stac @ {context} @ {thm rv2_Cons} 1 ⟩)
apply (tactic ⟨ stac @ {context} @ {thm sum.case(1)} 1 ⟩)
apply (tactic ⟨ stac @ {context} @ {thm frombd_F22_tobd_F22} 1 ⟩)
  apply assumption
apply (rule conjI)
apply (rule impI)
apply (rule conjI)
  apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic ⟨ dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 ⟩)
apply (drule mp)
  apply assumption
apply (tactic ⟨ dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 ⟩)
apply (drule mp)
  apply assumption
apply (tactic ⟨ dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 ⟩)
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI1)

```

```

apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (rule conjI)
apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @ {context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply assumption

```

```

apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (tactic << stac @ {context} @ {thm rv2_Cons} 1 >>)

```

```

apply (tactic << stac @{context} @{thm sum.case(2)} 1 >>)
apply (tactic << stac @{context} @{thm frombd_F23_tobd_F23} 1 >>)
apply assumption
apply (rule conjI)
apply (rule impI)
apply (rule conjI)
apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply assumption

apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply assumption

apply (rule impI)
apply (rule conjI)
apply (rule impI)
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)

```

```

apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>))
apply (erule mp)
apply assumption

```

```

apply (rule impI)
apply (rule set_mp[OF equalityD2])
apply (rule Lev2_Suc)
apply (rule ssubst_mem[OF append_Cons])
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply (erule conjI)
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (erule mp)
apply assumption
done

```

```

lemmas Fset_Lev' = spec[OF spec[OF spec[OF spec[OF spec[OF spec[OF spec[OF spec[OF Fset_Lev]]]]]]]]]
lemmas F1set2_Lev1 = mp[OF conjunct1[OF mp[OF conjunct1[OF mp[OF conjunct1[OF Fset_Lev]]]]]]]
lemmas F1set2_Lev2 = mp[OF conjunct1[OF mp[OF conjunct1[OF mp[OF conjunct2[OF Fset_Lev]]]]]]]
lemmas F2set2_Lev1 = mp[OF conjunct1[OF mp[OF conjunct2[OF mp[OF conjunct1[OF Fset_Lev]]]]]]]
lemmas F2set2_Lev2 = mp[OF conjunct1[OF mp[OF conjunct2[OF mp[OF conjunct2[OF Fset_Lev]]]]]]]
lemmas F1set3_Lev1 = mp[OF conjunct2[OF mp[OF conjunct1[OF mp[OF conjunct1[OF Fset_Lev]]]]]]]
lemmas F1set3_Lev2 = mp[OF conjunct2[OF mp[OF conjunct1[OF mp[OF conjunct2[OF Fset_Lev]]]]]]]
lemmas F2set3_Lev1 = mp[OF conjunct2[OF mp[OF conjunct2[OF mp[OF conjunct1[OF Fset_Lev]]]]]]]
lemmas F2set3_Lev2 = mp[OF conjunct2[OF mp[OF conjunct2[OF mp[OF conjunct2[OF Fset_Lev]]]]]]]

```

```

lemma Fset_image_Lev:
   $\forall kl\ k\ b1\ b2\ b1'\ b2'.$ 
  ( $kl \in Lev1\ s1\ s2\ n\ b1 \rightarrow$ 
    ( $kl @ [Inl\ k] \in Lev1\ s1\ s2\ (Suc\ n)\ b1 \rightarrow$ 
      ( $rv1\ s1\ s2\ kl\ b1 = Inl\ b1' \rightarrow k \in tobd\_F12\ s1\ b1' \text{ ' } F1set2\ (s1\ b1')$ )  $\wedge$ 
      ( $rv1\ s1\ s2\ kl\ b1 = Inr\ b2' \rightarrow k \in tobd\_F22\ s2\ b2' \text{ ' } F2set2\ (s2\ b2')$ ))  $\wedge$ 
      ( $kl @ [Inr\ k] \in Lev1\ s1\ s2\ (Suc\ n)\ b1 \rightarrow$ 
        ( $rv1\ s1\ s2\ kl\ b1 = Inl\ b1' \rightarrow k \in tobd\_F13\ s1\ b1' \text{ ' } F1set3\ (s1\ b1')$ )  $\wedge$ 
        ( $rv1\ s1\ s2\ kl\ b1 = Inr\ b2' \rightarrow k \in tobd\_F23\ s2\ b2' \text{ ' } F2set3\ (s2\ b2')$ )))))  $\wedge$ 
    ( $kl \in Lev2\ s1\ s2\ n\ b2 \rightarrow$ 
      ( $kl @ [Inl\ k] \in Lev2\ s1\ s2\ (Suc\ n)\ b2 \rightarrow$ 
        ( $rv2\ s1\ s2\ kl\ b2 = Inl\ b1' \rightarrow k \in tobd\_F12\ s1\ b1' \text{ ' } F1set2\ (s1\ b1')$ )  $\wedge$ 
        ( $rv2\ s1\ s2\ kl\ b2 = Inr\ b2' \rightarrow k \in tobd\_F22\ s2\ b2' \text{ ' } F2set2\ (s2\ b2')$ ))  $\wedge$ 
        ( $kl @ [Inr\ k] \in Lev2\ s1\ s2\ (Suc\ n)\ b2 \rightarrow$ 
          ( $rv2\ s1\ s2\ kl\ b2 = Inl\ b1' \rightarrow k \in tobd\_F13\ s1\ b1' \text{ ' } F1set3\ (s1\ b1')$ )  $\wedge$ 
          ( $rv2\ s1\ s2\ kl\ b2 = Inr\ b2' \rightarrow k \in tobd\_F23\ s2\ b2' \text{ ' } F2set3\ (s2\ b2')$ )))))
    )
apply (rule nat_induct[of _ n])

```

```

apply (rule allI)+
apply (rule conjI)
apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev1_0]])
apply (erule singletonE)

```

```

apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (rule conjI)
apply (rule impI)
apply (drule trans[OF sym])
  apply (rule rv1_Nil)
apply (drule ssubst_mem[OF sym[OF append_Nil]])
apply (drule set_mp[OF equalityD1[OF Lev1_Suc]])
apply (drule Inl_inject)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (erule UnE)
  apply (erule CollectE exE conjE)+
  apply (drule list.inject[THEN iffD1])
  apply (erule conjE)
  apply (drule Inl_inject)
  apply (tactic << hyp_subst_tac @{context} 1 >>)
  apply (erule imageI)
apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (erule notE[OF Inl_not_Inr])
apply (rule impI)
apply (drule trans[OF sym])
apply (rule rv1_Nil)
apply (erule notE[OF Inr_not_Inl])

```

```

apply (rule impI)
apply (rule conjI)
apply (rule impI)
apply (drule ssubst_mem[OF sym[OF append_Nil]])
apply (drule set_mp[OF equalityD1[OF Lev1_Suc]])
apply (drule trans[OF sym])
  apply (rule rv1_Nil)
apply (drule Inl_inject)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (erule UnE)
  apply (erule CollectE exE conjE)+
  apply (drule list.inject[THEN iffD1])
  apply (erule conjE)
  apply (erule notE[OF Inr_not_Inl])
apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (drule Inr_inject)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (erule imageI)
apply (rule impI)
apply (drule trans[OF sym])
  apply (rule rv1_Nil)
apply (erule notE[OF Inr_not_Inl])

```

```

apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev2_0]])
apply (erule singletonE)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (rule conjI)
  apply (rule impI)
  apply (drule trans[OF sym])
  apply (rule rv2_Nil)
apply (erule notE[OF Inl_not_Inr])

```



```

apply (rule impI)
apply (drule ssubst_mem[OF sym[OF append_Nil]])
apply (drule set_mp[OF equalityD1[OF Lev2_Suc]])
apply (drule trans[OF sym])
  apply (rule rv2_Nil)
apply (drule Inr_inject)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (erule UnE)
  apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (drule Inl_inject)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (erule imageI)
apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (erule notE[OF Inl_not_Inr])

```

```

apply (rule impI)
apply (rule conjI)
apply (rule impI)
apply (drule trans[OF sym])
  apply (rule rv2_Nil)
apply (erule notE[OF Inl_not_Inr])
apply (rule impI)
apply (drule ssubst_mem[OF sym[OF append_Nil]])
apply (drule set_mp[OF equalityD1[OF Lev2_Suc]])
apply (drule trans[OF sym])
  apply (rule rv2_Nil)
apply (drule Inr_inject)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (erule UnE)
  apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (erule notE[OF Inr_not_Inl])
apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (drule Inr_inject)
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (erule imageI)

```

```

apply (rule allI)+
apply (rule conjI)
apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev1_Suc]])
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule conjI)
  apply (rule impI)
  apply (drule ssubst_mem[OF sym[OF append_Cons]])
  apply (drule set_mp[OF equalityD1[OF Lev1_Suc]])
  apply (erule UnE)
  apply (erule CollectE exE conjE)+
  apply (drule list.inject[THEN iffD1])
  apply (erule conjE)
  apply (drule Inl_inject)
  apply (tactic << dtac @{context}

```

```

(Thm.permute_prem 0 2 (@{thm tobd_F12_inj} RS iffD1)) 1 ))
  apply assumption
  apply assumption
  apply (tactic << hyp_subst_tac @{context} 1 >>)
  apply (rule conjI)
  apply (rule impI)
  apply (drule trans[OF sym])
  apply (rule trans[OF rv1_Cons])
  apply (rule trans[OF arg_cong[OF sum.case(1)]])
  apply (erule arg_cong[OF frombd_F12_tobd_F12])
  apply (erule allE)+
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
  apply (drule mp)
  apply assumption
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
  apply (drule mp)
  apply assumption
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
  apply (erule mp)
  apply (erule sym)

  apply (rule impI)
  apply (drule trans[OF sym])
  apply (rule trans[OF rv1_Cons])
  apply (rule trans[OF arg_cong[OF sum.case(1)]])
  apply (erule arg_cong[OF frombd_F12_tobd_F12])
  apply (erule allE)+
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
  apply (drule mp)
  apply assumption
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
  apply (drule mp)
  apply assumption
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
  apply (erule mp)
  apply (erule sym)

  apply (erule CollectE exE conjE)+
  apply (drule list.inject[THEN iffD1])
  apply (erule conjE)
  apply (erule notE[OF Inl_not_Inr])

  apply (rule impI)
  apply (drule ssubst_mem[OF sym[OF append_Cons]])
  apply (drule set_mp[OF equalityD1[OF Lev1_Suc]])
  apply (erule UnE)
  apply (erule CollectE exE conjE)+
  apply (drule list.inject[THEN iffD1])
  apply (erule conjE)
  apply (drule Inl_inject)
  apply (tactic << dtac @{context}
(Thm.permute_prem 0 2 (@{thm tobd_F12_inj} THEN iffD1)) 1 ))
  apply assumption
  apply assumption
  apply (tactic << hyp_subst_tac @{context} 1 >>)
  apply (rule conjI)
  apply (rule impI)
  apply (drule trans[OF sym])
  apply (rule trans[OF rv1_Cons])
  apply (rule trans[OF arg_cong[OF sum.case(1)]])
  apply (erule arg_cong[OF frombd_F12_tobd_F12])
  apply (erule allE)+
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)

```

```

apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply (erule sym)

apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF arg_cong[OF sum.case(1)]])
apply (erule arg_cong[OF frombd_F12_tobd_F12])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply (erule sym)

apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (erule notE[OF Inl_not_Inr])

apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (drule ssubst_mem[OF sym[OF append_Cons]])
apply (drule set_mp[OF equalityD1[OF Lev1_Suc]])
apply (erule UnE)
  apply (erule CollectE exE conjE)+
  apply (drule list.inject[THEN iffD1])
  apply (erule conjE)
  apply (erule notE[OF Inr_not_Inl])

apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (drule Inr_inject)
apply (tactic << dtac @{context}
(Thm.permute_premis 0 2 @{thm tobd_F13_inj[THEN iffD1]}) 1 >>)
  apply assumption
  apply assumption
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply (erule arg_cong[OF frombd_F13_tobd_F13])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)

```

```

apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>))
apply (erule mp)
apply (erule sym)

apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply (erule arg_cong[OF frombd_F13_tobd_F13])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>))
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (erule mp)
apply (erule sym)

apply (rule impI)
apply (drule subst_mem[OF sym[OF append_Cons]])
apply (drule set_mp[OF equalityD1[OF Lev1_Suc]])
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (erule notE[OF Inr_not_Inl])

apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (drule Inr_inject)
apply (tactic << dtac @{context}
(Thm.permute_prems 0 2 @{thm tobd_F13_inj[THEN iffD1]} 1 >>))
apply assumption
apply assumption
apply (tactic << hyp_subst_tac @{context} 1 >>))
apply (rule conjI)
apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply (erule arg_cong[OF frombd_F13_tobd_F13])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>))
apply (erule mp)
apply (erule sym)

apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply (erule arg_cong[OF frombd_F13_tobd_F13])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>))

```

```

apply (drule mp)
apply assumption
apply (tactic << dtac @{\context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @{\context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply (erule sym)

```

```

apply (rule impI)
apply (drule set_mp[OF equalityD1[OF Lev2_Suc]])
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @{\context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (drule ssubst_mem[OF sym[OF append_Cons]])
apply (drule set_mp[OF equalityD1[OF Lev2_Suc]])
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (drule Inl_inject)
apply (tactic << dtac @{\context}
(Thm.permute_premis 0 2 @{\thm tobd_F22_inj[THEN iffD1]}) 1 >>)
apply assumption
apply assumption
apply (tactic << hyp_subst_tac @{\context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv2_Cons])
apply (rule trans[OF arg_cong[OF sum.case(1)]])
apply (erule arg_cong[OF frombd_F22_tobd_F22])
apply (erule allE)+
apply (tactic << dtac @{\context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @{\context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @{\context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply (erule sym)

```

```

apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv2_Cons])
apply (rule trans[OF arg_cong[OF sum.case(1)]])
apply (erule arg_cong[OF frombd_F22_tobd_F22])
apply (erule allE)+
apply (tactic << dtac @{\context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @{\context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @{\context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply (erule sym)

```

```

apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (erule notE[OF Inl_not_Inr])

apply (rule impI)
apply (drule ssubst_mem[OF sym[OF append_Cons]])
apply (drule set_mp[OF equalityD1[OF Lev2_Suc]])
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (drule Inl_inject)
apply (tactic << dtac @{context}
(Thm.permute_premis 0 2 @{thm tobd_F22_inj[THEN iffD1]}) 1 >>)
  apply assumption
  apply assumption
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv2_Cons])
apply (rule trans[OF arg_cong[OF sum.case(1)]])
apply (erule arg_cong[OF frombd_F22_tobd_F22])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply (erule sym)

apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv2_Cons])
apply (rule trans[OF arg_cong[OF sum.case(1)]])
apply (erule arg_cong[OF frombd_F22_tobd_F22])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply (erule sym)

apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (erule notE[OF Inl_not_Inr])

apply (erule CollectE exE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (drule ssubst_mem[OF sym[OF append_Cons]])
apply (drule set_mp[OF equalityD1[OF Lev2_Suc]])
apply (erule UnE)

```

```

apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (erule notE[OF Inr_not_Inl])

apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (drule Inr_inject)
apply (tactic << dtac @{context}
(Thm.permute_premis 0 2 @{thm tobd_F23_inj[THEN iffD1]}) 1 >>)
  apply assumption
apply assumption
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (drule trans[OF sym])
  apply (rule trans[OF rv2_Cons])
  apply (rule trans[OF arg_cong[OF sum.case(2)]])
  apply (erule arg_cong[OF frombd_F23_tobd_F23])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply (erule sym)

apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv2_Cons])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply (erule arg_cong[OF frombd_F23_tobd_F23])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (drule mp)
  apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply (erule sym)

apply (rule impI)
apply (drule ssubst_mem[OF sym[OF append_Cons]])
apply (drule set_mp[OF equalityD1[OF Lev2_Suc]])
apply (erule UnE)
  apply (erule CollectE exE conjE)+
  apply (drule list.inject[THEN iffD1])
  apply (erule conjE)
  apply (erule notE[OF Inr_not_Inl])

apply (erule CollectE exE conjE)+
apply (drule list.inject[THEN iffD1])
apply (erule conjE)
apply (drule Inr_inject)
apply (tactic << dtac @{context}
(Thm.permute_premis 0 2 @{thm tobd_F23_inj[THEN iffD1]}) 1 >>)
  apply assumption

```

```

apply assumption
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule conjI)
apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv2_Cons])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply (erule arg_cong[OF frombd_F23_tobd_F23])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (erule mp)
apply (erule sym)

```

```

apply (rule impI)
apply (drule trans[OF sym])
apply (rule trans[OF rv2_Cons])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply (erule arg_cong[OF frombd_F23_tobd_F23])
apply (erule allE)+
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (drule mp)
apply assumption
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply (erule sym)
done

```

```

lemmas Fset_image_Lev' =
  spec[OF spec[OF spec[OF spec[OF spec[OF spec[OF Fset_image_Lev]]]]]]]
lemmas F1set2_image_Lev1 =
  mp[OF conjunct1[OF mp[OF conjunct1[OF mp[OF conjunct1[OF Fset_image_Lev']]]]]]]
lemmas F1set2_image_Lev2 =
  mp[OF conjunct1[OF mp[OF conjunct1[OF mp[OF conjunct2[OF Fset_image_Lev']]]]]]]
lemmas F1set3_image_Lev1 =
  mp[OF conjunct1[OF mp[OF conjunct2[OF mp[OF conjunct1[OF Fset_image_Lev']]]]]]]
lemmas F1set3_image_Lev2 =
  mp[OF conjunct1[OF mp[OF conjunct2[OF mp[OF conjunct2[OF Fset_image_Lev']]]]]]]
lemmas F2set2_image_Lev1 =
  mp[OF conjunct2[OF mp[OF conjunct1[OF mp[OF conjunct1[OF Fset_image_Lev']]]]]]]
lemmas F2set2_image_Lev2 =
  mp[OF conjunct2[OF mp[OF conjunct1[OF mp[OF conjunct2[OF Fset_image_Lev']]]]]]]
lemmas F2set3_image_Lev1 =
  mp[OF conjunct2[OF mp[OF conjunct2[OF mp[OF conjunct1[OF Fset_image_Lev']]]]]]]
lemmas F2set3_image_Lev2 =
  mp[OF conjunct2[OF mp[OF conjunct2[OF mp[OF conjunct2[OF Fset_image_Lev']]]]]]]

```

```

lemma mor_beh:
  mor UNIV UNIV s1 s2 carT1 carT2 strT1 strT2 (beh1 s1 s2) (beh2 s1 s2)
apply (rule mor_cong)
apply (rule ext[OF beh1_def])
apply (rule ext[OF beh2_def])
apply (tactic << rtac @{context} (@{thm mor_def} RS iffD2) 1 >>)
apply (rule conjI)
apply (rule conjI)

```



```

apply (rule ballI)
apply (rule set_mp[OF equalityD2[OF carT1_def]])
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (rule conjI)
apply (rule conjI)
  apply (rule UN_I)
    apply (rule UNIV_I)
      apply (rule set_mp)
        apply (rule equalityD2)
          apply (rule Lev1_0)
            apply (rule singletonI)

apply (rule ballI)
apply (erule UN_E)
apply (rule conjI)
apply (rule ballI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (rule rev_mp[OF rv1_Inl_last_impI])
apply (erule exE)
apply (rule iffD2[OF isNode1_def])
apply (rule exI)
apply (rule conjI)
  apply (erule trans[OF sum.case_cong_weak])
  apply (rule sum.case(1))

apply (rule conjI)
apply (rule trans[OF F1.set_map(2)])
apply (rule equalityI)
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule SuccI)
  apply (rule UN_I[OF UNIV_I])
  apply (erule F1set2_Lev1)
  apply assumption
  apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule F1set2_image_Lev1)
  apply assumption
apply (drule length_Lev1)
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (drule length_Lev1)
apply (erule set_mp[OF equalityD1[OF arg_cong[OF length_append_singleton]]])
apply assumption

apply (rule trans[OF F1.set_map(3)])
apply (rule equalityI)
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule SuccI)
  apply (rule UN_I[OF UNIV_I])
  apply (erule F1set3_Lev1)
  apply assumption
  apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+

```

```

apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule F1set3_image_Lev1)
  apply assumption
  apply (drule length_Lev1)
  apply (tactic ⟨⟨ hyp_subst_tac @ {context} 1 ⟩⟩)
  apply (drule length_Lev1)
  apply (erule set_mp[OF equalityD1[OF arg_cong[OF length_append_singleton]]])
apply assumption

```

```

apply (rule ballI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (rule rev_mp[OF rv1_Inr_last_impI])
apply (erule exE)
apply (rule iffD2[OF isNode2_def])
apply (rule exI)
apply (rule conjI)
apply (erule trans[OF sum.case_cong_weak])
apply (rule sum.case(2))

```

```

apply (rule conjI)
apply (rule trans[OF F2.set_map(2)])
apply (rule equalityI)
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule SuccI)
  apply (rule UN_I[OF UNIV_I])
  apply (erule F2set2_Lev1)
  apply assumption
  apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule F2set2_image_Lev1)
  apply assumption
  apply (drule length_Lev1)
  apply (tactic ⟨⟨ hyp_subst_tac @ {context} 1 ⟩⟩)
  apply (drule length_Lev1)
  apply (erule set_mp[OF equalityD1[OF arg_cong[OF length_append_singleton]]])
apply assumption

```

```

apply (rule trans[OF F2.set_map(3)])
apply (rule equalityI)
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule SuccI)
  apply (rule UN_I[OF UNIV_I])
  apply (erule F2set3_Lev1)
  apply assumption
  apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule F2set3_image_Lev1)
  apply assumption

```

```

apply (drule length_Lev1)
apply (tactic (<< hyp_subst_tac @ {context} 1 >>))
apply (drule length_Lev1')
apply (erule set_mp[OF equalityD1[OF arg_cong[OF length_append_singleton]]])
apply assumption

```

```

apply (rule iffD2[OF isNode1_def])
apply (rule exI)
apply (rule conjI)
apply (rule trans[OF sum.case_cong_weak])
apply (rule rv1_Nil)
apply (rule sum.case(1))

```

```

apply (rule conjI)
apply (rule trans[OF F1.set_map(2)])
apply (rule equalityI)
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule SuccI)
apply (rule UN_I[OF UNIV_I])
apply (rule F1set2_Lev1)
apply (rule set_mp[OF equalityD2])
apply (rule Lev1_0)
apply (rule singletonI)
apply (rule rv1_Nil)
apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (rule F1set2_image_Lev1)
apply (rule set_mp[OF equalityD2[OF Lev1_0]])
apply (rule singletonI)
apply (drule length_Lev1')
apply (erule set_mp[OF equalityD1[OF arg_cong[OF
  trans[OF length_append_singleton arg_cong[of _ _ Suc, OF list.size(3)]]]])
apply (rule rv1_Nil)

```

```

apply (rule trans[OF F1.set_map(3)])
apply (rule equalityI)
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule SuccI)
apply (rule UN_I[OF UNIV_I])
apply (rule F1set3_Lev1)
apply (rule set_mp[OF equalityD2])
apply (rule Lev1_0)
apply (rule singletonI)
apply (rule rv1_Nil)
apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (rule F1set3_image_Lev1)
apply (rule set_mp[OF equalityD2[OF Lev1_0]])
apply (rule singletonI)
apply (drule length_Lev1')
apply (erule set_mp[OF equalityD1[OF arg_cong[OF
  trans[OF length_append_singleton arg_cong[of _ _ Suc, OF list.size(3)]]]])
apply (rule rv1_Nil)

```

```

apply (rule ballI)
apply (rule set_mp[OF equalityD2[OF carT2_def]])
apply (rule CollectI)

```

```

apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply (rule conjI)
apply (rule conjI)
apply (rule UN_I)
  apply (rule UNIV_I)
apply (rule set_mp)
  apply (rule equalityD2)
  apply (rule Lev2_0)
apply (rule singletonI)

apply (rule ballI)
apply (erule UN_E)
apply (rule conjI)
apply (rule ballI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (rule rev_mp[OF rv2_Inl_last_impI])
apply (erule exE)
apply (rule iffD2[OF isNode1_def])
apply (rule exI)
apply (rule conjI)
apply (erule trans[OF sum.case_cong_weak])
apply (rule sum.case(1))

apply (rule conjI)
apply (rule trans[OF F1.set_map(2)])
apply (rule equalityI)
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule SuccI)
  apply (rule UN_I[OF UNIV_I])
  apply (erule F1set2_Lev2)
  apply assumption
apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule F1set2_image_Lev2)
  apply assumption
  apply (drule length_Lev2)
  apply (tactic ⟨ hyp_subst_tac @ {context} 1 ⟩)
  apply (drule length_Lev2')
apply (erule set_mp[OF equalityD1[OF arg_cong[OF length_append_singleton]])
apply assumption

apply (rule trans[OF F1.set_map(3)])
apply (rule equalityI)
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule SuccI)
apply (rule UN_I[OF UNIV_I])
apply (erule F1set3_Lev2)
  apply assumption
apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)

```

```

apply (erule thin_rl)
apply (rule F1set3_image_Lev2)
  apply assumption
  apply (drule length_Lev2)
  apply (tactic << hyp_subst_tac @ {context} 1 >>)
  apply (drule length_Lev21)
  apply (erule set_mp[OF equalityD1[OF arg_cong[OF length_append_singleton]]])
apply assumption

```

```

apply (rule ballI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (rule rev_mp[OF rv2_Inr_last_impI])
apply (erule exE)
apply (rule iffD2[OF isNode2_def])
apply (rule exI)
apply (rule conjI)
  apply (erule trans[OF sum.case_cong_weak])
  apply (rule sum.case(2))

```

```

apply (rule conjI)
apply (rule trans[OF F2.set_map(2)])
apply (rule equalityI)
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule SuccI)
  apply (rule UN_I[OF UNIV_I])
  apply (erule F2set2_Lev2)
  apply assumption
  apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule F2set2_image_Lev2)
  apply assumption
  apply (drule length_Lev2)
  apply (tactic << hyp_subst_tac @ {context} 1 >>)
  apply (drule length_Lev21)
  apply (erule set_mp[OF equalityD1[OF arg_cong[OF length_append_singleton]]])
apply assumption

```

```

apply (rule trans[OF F2.set_map(3)])
apply (rule equalityI)
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule SuccI)
  apply (rule UN_I[OF UNIV_I])
  apply (erule F2set3_Lev2)
  apply assumption
  apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule F2set3_image_Lev2)
  apply assumption
  apply (drule length_Lev2)
  apply (tactic << hyp_subst_tac @ {context} 1 >>)
  apply (drule length_Lev21)

```

```

apply (erule set_mp[OF equalityD1[OF arg_cong[OF length_append_singleton]]])
apply assumption

apply (rule iffD2[OF isNode2_def])
apply (rule exI)
apply (rule conjI)
apply (rule trans[OF sum.case_cong_weak])
  apply (rule rv2_Nil)
apply (rule sum.case(2))

apply (rule conjI)
apply (rule trans[OF F2.set_map(2)])
apply (rule equalityI)
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule SuccI)
  apply (rule UN_I[OF UNIV_I])
  apply (rule F2set2_Lev2)
    apply (rule set_mp[OF equalityD2[OF Lev2_0]])
    apply (rule singletonI)
    apply (rule rv2_Nil)
  apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (rule F2set2_image_Lev2)
  apply (rule set_mp[OF equalityD2[OF Lev2_0]])
  apply (rule singletonI)
apply (drule length_Lev2')
apply (erule set_mp[OF equalityD1[OF arg_cong[OF
  trans[OF length_append_singleton arg_cong[of _ _ Suc, OF list.size(3)]]]])
apply (rule rv2_Nil)

apply (rule trans[OF F2.set_map(3)])
apply (rule equalityI)
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule SuccI)
apply (rule UN_I[OF UNIV_I])
apply (rule F2set3_Lev2)
  apply (rule set_mp[OF equalityD2])
    apply (rule Lev2_0)
    apply (rule singletonI)
  apply (rule rv2_Nil)
apply assumption
apply (rule subsetI)
apply (erule CollectE SuccD[elim_format] UN_E)+
apply (rule F2set3_image_Lev2)
  apply (rule set_mp[OF equalityD2[OF Lev2_0]])
  apply (rule singletonI)
apply (drule length_Lev2')
apply (erule set_mp[OF equalityD1[OF arg_cong[OF
  trans[OF length_append_singleton arg_cong[of _ _ Suc, OF list.size(3)]]]])
apply (rule rv2_Nil)

apply (rule conjI)
apply (rule ballI)
apply (rule sym)
apply (rule trans)
  apply (rule trans[OF fun_cong[OF strT1_def] prod.case])
apply (tactic ⟨⟨ CONVERSION (Conv.top_conv
  (K (Conv.try_conv (Conv.rewr_conv (@{thm rv1_Nil} RS eq_reflection)))) @context 1 ⟩⟩)

```

```

apply (rule trans[OF sum.case_cong_weak])
apply (rule sum.case(1))
apply (rule trans[OF sum.case(1)])
apply (rule trans[OF F1map_comp_id])
apply (rule F1.map_cong0[OF refl])
apply (rule trans)
  apply (rule o_apply)
apply (rule iffD2)
  apply (rule prod.inject)
apply (rule conjI)
apply (rule trans)
  apply (rule Shift_def)

apply (rule equalityI)
apply (rule subsetI)
apply (erule thin_rl)
apply (erule CollectE UN_E)+
apply (erule length_Lev1')
apply (erule asm_rl)
apply (erule thin_rl)
apply (erule set_rev_mp[OF equalityD1])
  apply (rule trans[OF arg_cong[OF length_Cons]])
apply (rule Lev1_Suc)
apply (erule UnE)
  apply (erule CollectE exE conjE)+
  apply (tactic ⟨ dtac @context ⟩ @thm list.inject[THEN iffD1] 1 ⟩)
apply (erule conjE)
apply (erule Inl_inject)
apply (tactic ⟨ dtac @context ⟩
(Thm.permute_prem 0 2 @thm tobd_F12_inj[THEN iffD1]) 1 ⟩)
  apply assumption
  apply assumption
  apply (tactic ⟨ hyp_subst_tac @context ⟩ 1 ⟩)
apply (erule UN_I[OF UNIV_I])
apply (erule CollectE exE conjE)+
apply (tactic ⟨ dtac @context ⟩ @thm list.inject[THEN iffD1] 1 ⟩)
apply (erule conjE)
apply (erule notE[OF Inl_not_Inr])

apply (rule UN_least)
apply (rule subsetI)
apply (rule CollectI)
apply (rule UN_I[OF UNIV_I])
apply (rule set_mp[OF equalityD2])
  apply (rule Lev1_Suc)
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply assumption

apply (rule trans)
  apply (rule shift_def)
apply (rule iffD2)
  apply (rule fun_eq_iff)
apply (rule allI)
apply (tactic ⟨ CONVERSION (Conv.top_conv
(K (Conv.try_conv (Conv.rewr_conv (@thm rv1_Cons) RS eq_reflection)))) @context ⟩ 1 ⟩)
apply (rule sum.case_cong_weak)

```

```

apply (rule trans[OF sum.case(1)])
apply (drule frombd_F12_tobd_F12)
apply (erule arg_cong)

apply (rule trans)
apply (rule o_apply)
apply (rule iffD2)
apply (rule prod.inject)
apply (rule conjI)
apply (rule trans)
apply (rule Shift_def)

apply (rule equalityI)
apply (rule subsetI)
apply (erule thin_rl)
apply (erule CollectE UN_E)+
apply (drule length_Lev1⌈)
apply (drule asm_rl)
apply (erule thin_rl)
apply (drule set_rev_mp[OF _ equalityD1])
apply (rule trans[OF arg_cong[OF length_Cons]])
apply (rule Lev1_Suc)
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (tactic ⟨ dtac @ {context} @ {thm list.inject[THEN iffD1]} 1 ⟩)
apply (erule conjE)
apply (erule notE[OF Inr_not_Inl])
apply (erule CollectE exE conjE)+
apply (tactic ⟨ dtac @ {context} @ {thm list.inject[THEN iffD1]} 1 ⟩)
apply (erule conjE)
apply (drule Inr_inject)
apply (tactic ⟨ dtac @ {context}
(Thm.permute_premis 0 2 @ {thm tobd_F13_inj[THEN iffD1]}) 1 ⟩)
apply assumption
apply assumption
apply (tactic ⟨ hyp_subst_tac @ {context} 1 ⟩)
apply (erule UN_I[OF UNIV_I])

apply (rule UN_least)
apply (rule subsetI)
apply (rule CollectI)
apply (rule UN_I[OF UNIV_I])
apply (rule set_mp[OF equalityD2])
apply (rule Lev1_Suc)
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply (erule conjI)
apply assumption

apply (rule trans)
apply (rule shift_def)
apply (rule iffD2)
apply (rule fun_eq_iff)
apply (rule allI)
apply (rule sum.case_cong_weak)
apply (rule trans[OF rv1_Cons])
apply (rule trans[OF sum.case(2)])
apply (erule arg_cong[OF frombd_F13_tobd_F13])

```



```

apply (rule ballI)
apply (rule sym)
apply (rule trans)
  apply (rule trans[OF fun_cong[OF strT2_def] prod.case])
apply (rule trans[OF sum.case_cong_weak[OF trans[OF sum.case_cong_weak]]])
  apply (rule rv2_Nil)
  apply (rule sum.case(2))
apply (rule trans[OF sum.case(2)])
apply (rule trans[OF F2map_comp_id])
apply (rule F2.map_cong0[OF refl])
apply (rule trans)
  apply (rule o_apply)
apply (rule iffD2)
apply (rule prod.inject)
apply (rule conjI)
apply (rule trans)
  apply (rule Shift_def)

```

```

apply (rule equalityI)
apply (rule subsetI)
apply (erule thin_rl)
apply (erule CollectE UN_E)+
apply (erule length_Lev2⌈)
apply (erule asm_rl)
apply (erule thin_rl)
apply (erule set_rev_mp[OF _equalityD1])
  apply (rule trans[OF arg_cong[OF length_Cons]])
  apply (rule Lev2_Suc)
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (tactic << dtac @{context} @{thm list.inject[THEN iffD1]} 1 >>)
apply (erule conjE)
apply (erule Inl_inject)
apply (tactic << dtac @{context}
(Thm.permute_premis 0 2 @{thm tobd_F22_inj[THEN iffD1]} 1 >>)
  apply assumption
  apply assumption
  apply (tactic << hyp_subst_tac @{context} 1 >>)
  apply (erule UN_I[OF UNIV_I])
apply (erule CollectE exE conjE)+
apply (tactic << dtac @{context} @{thm list.inject[THEN iffD1]} 1 >>)
apply (erule conjE)
apply (erule notE[OF Inl_not_Inr])

```

```

apply (rule UN_least)
apply (rule subsetI)
apply (rule CollectI)
apply (rule UN_I[OF UNIV_I])
apply (rule set_mp[OF equalityD2])
  apply (rule Lev2_Suc)
apply (rule UnI1)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
  apply (rule refl)
apply (erule conjI)
apply assumption

```

```

apply (rule trans)
apply (rule shift_def)
apply (rule iffD2)
apply (rule fun_eq_iff)
apply (rule allI)
apply (rule sum.case_cong_weak)
apply (rule trans[OF rv2_Cons])
apply (rule trans[OF arg_cong[OF sum.case(1)]])
apply (erule arg_cong[OF frombd_F22_tobd_F22])

```

```

apply (rule trans)
apply (rule o_apply)
apply (rule iffD2)
apply (rule prod.inject)
apply (rule conjI)
apply (rule trans)
apply (rule Shift_def)

```

```

apply (rule equalityI)
apply (rule subsetI)
apply (erule thin_rl)
apply (erule CollectE UN_E)+
apply (drule length_Lev2')
apply (drule asm_rl)
apply (erule thin_rl)
apply (drule set_rev_mp[OF equalityD1])
apply (rule trans[OF arg_cong[OF length_Cons]])
apply (rule Lev2_Suc)
apply (erule UnE)
apply (erule CollectE exE conjE)+
apply (tactic << dtac @ {context} @ {thm list.inject[THEN iffD1]} 1 >>)
apply (erule conjE)
apply (erule notE[OF Inr_not_Inl])
apply (erule CollectE exE conjE)+
apply (tactic << dtac @ {context} @ {thm list.inject[THEN iffD1]} 1 >>)
apply (erule conjE)
apply (drule Inr_inject)
apply (tactic << dtac @ {context}
(Thm.permute_premis 0 2 @ {thm tobd_F23_inj[THEN iffD1]} 1 >>)
apply assumption
apply assumption
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (erule UN_I[OF UNIV_I])

```

```

apply (rule UN_least)
apply (rule subsetI)
apply (rule CollectI)
apply (rule UN_I[OF UNIV_I])
apply (rule set_mp[OF equalityD2])
apply (rule Lev2_Suc)
apply (rule UnI2)
apply (rule CollectI)
apply (rule exI)+
apply (rule conjI)
apply (rule refl)
apply (erule conjI)
apply assumption

```

```

apply (rule trans)
apply (rule shift_def)
apply (rule iffD2)

```

apply (rule fun_eq_iff)
apply (rule allI)

apply (rule sum.case_cong_weak)
apply (rule trans[OF rv2_Cons])
apply (rule trans[OF arg_cong[OF sum.case(2)]])
apply (erule arg_cong[OF frombd_F23_tobd_F23])
done

2.6 Quotient Coalgebra

abbreviation car_final1 **where**
 car_final1 \equiv carT1 // (lsbis1 carT1 carT2 strT1 strT2)

abbreviation car_final2 **where**
 car_final2 \equiv carT2 // (lsbis2 carT1 carT2 strT1 strT2)

abbreviation str_final1 **where**
 str_final1 \equiv univ (F1map id
 (Equiv_Relations.proj (lsbis1 carT1 carT2 strT1 strT2))
 (Equiv_Relations.proj (lsbis2 carT1 carT2 strT1 strT2))) o strT1)

abbreviation str_final2 **where**
 str_final2 \equiv univ (F2map id
 (Equiv_Relations.proj (lsbis1 carT1 carT2 strT1 strT2))
 (Equiv_Relations.proj (lsbis2 carT1 carT2 strT1 strT2))) o strT2)

lemma congruent_strQ1: congruent (lsbis1 carT1 carT2 strT1 strT2 :: 'a carrier rel)
 (F1map id (Equiv_Relations.proj (lsbis1 carT1 carT2 strT1 strT2 :: 'a carrier rel))
 (Equiv_Relations.proj (lsbis2 carT1 carT2 strT1 strT2 :: 'a carrier rel))) o strT1)

apply (rule congruentI)
apply (erule lsbisE1)
apply (erule bezE conjE CollectE)+
apply (rule trans[OF o_apply])
apply (erule trans[OF arg_cong[OF sym]])
apply (rule trans[OF F1map_comp_id])
apply (rule trans[OF F1.map_cong0])
apply (rule refl)
apply (rule equiv_proj)
apply (rule equiv_lsbis1)
apply (rule coalg_T)
apply (erule set_mp)
apply assumption
apply (rule equiv_proj)
apply (rule equiv_lsbis2)
apply (rule coalg_T)
apply (erule set_mp)
apply assumption
apply (rule sym)
apply (rule trans[OF o_apply])
apply (erule trans[OF arg_cong[OF sym]])
apply (rule F1map_comp_id)
done

lemma congruent_strQ2: congruent (lsbis2 carT1 carT2 strT1 strT2 :: 'a carrier rel)
 (F2map id (Equiv_Relations.proj (lsbis1 carT1 carT2 strT1 strT2 :: 'a carrier rel))
 (Equiv_Relations.proj (lsbis2 carT1 carT2 strT1 strT2 :: 'a carrier rel))) o strT2)

apply (rule congruentI)
apply (erule lsbisE2)
apply (erule bezE conjE CollectE)+
apply (rule trans[OF o_apply])
apply (erule trans[OF arg_cong[OF sym]])
apply (rule trans[OF F2map_comp_id])
apply (rule trans[OF F2.map_cong0])
apply (rule refl)
apply (rule equiv_proj)
apply (rule equiv_lsbis1)

```

  apply (rule coalg_T)
  apply (erule set_mp)
  apply assumption
  apply (rule equiv_proj)
  apply (rule equiv_lsbis2)
  apply (rule coalg_T)
  apply (erule set_mp)
  apply assumption
  apply (rule sym)
  apply (rule trans[OF o_apply])
  apply (erule trans[OF arg_cong[OF sym]])
  apply (rule F2map_comp_id)
done

```

lemma *coalg_final*:

```

  coalg car_final1 car_final2 str_final1 str_final2
  apply (tactic << rtac @{context} (@{thm coalg_def} RS iffD2) 1 >>)
  apply (rule conjI)
  apply (rule univ_preserves)
    apply (rule equiv_lsbis1)
    apply (rule coalg_T)
  apply (rule congruent_strQ1)
  apply (rule ballI)
  apply (rule ssubst_mem)
  apply (rule o_apply)
  apply (rule CollectI)
  apply (rule conjI)
  apply (rule subset_UNIV)
  apply (rule conjI)
  apply (rule ord_eq_le_trans[OF F1.set_map(2)])
  apply (rule image_subsetI)
  apply (rule iffD2)
    apply (rule proj_in_iff)
    apply (rule equiv_lsbis1[OF coalg_T])
  apply (erule set_rev_mp)
  apply (erule coalg_F1set2[OF coalg_T])
  apply (rule ord_eq_le_trans[OF F1.set_map(3)])
  apply (rule image_subsetI)
  apply (rule iffD2)
    apply (rule proj_in_iff)
    apply (rule equiv_lsbis2[OF coalg_T])
  apply (erule set_rev_mp)
  apply (erule coalg_F1set3[OF coalg_T])

```

```

  apply (rule univ_preserves)
    apply (rule equiv_lsbis2)
    apply (rule coalg_T)
  apply (rule congruent_strQ2)
  apply (rule ballI)
  apply (tactic << stac @{context} @{thm o_apply} 1 >>)
  apply (rule CollectI)
  apply (rule conjI)
  apply (rule subset_UNIV)
  apply (rule conjI)
  apply (rule ord_eq_le_trans[OF F2.set_map(2)])
  apply (rule image_subsetI)
  apply (rule iffD2)
    apply (rule proj_in_iff)
    apply (rule equiv_lsbis1[OF coalg_T])
  apply (erule set_rev_mp)
  apply (erule coalg_F2set2[OF coalg_T])
  apply (rule ord_eq_le_trans[OF F2.set_map(3)])
  apply (rule image_subsetI)

```

```

apply (rule iffD2)
apply (rule proj_in_iff)
apply (rule equiv_lsbis2[OF coalg_T])
apply (erule set_rev_mp)
apply (erule coalg_F2set3[OF coalg_T])
done

```

lemma *mor_T_final*:

```

mor carT1 carT2 strT1 strT2 car_final1 car_final2 str_final1 str_final2
(Equiv_Relations.proj (lsbis1 carT1 carT2 strT1 strT2))
(Equiv_Relations.proj (lsbis2 carT1 carT2 strT1 strT2))
apply (tactic << rtac @{context} (@{thm mor_def} RS iffD2) 1 >>)
apply (rule conjI)
apply (rule conjI)
apply (rule ballI)
apply (rule iffD2)
apply (rule proj_in_iff)
apply (rule equiv_lsbis1[OF coalg_T])
apply assumption
apply (rule ballI)
apply (rule iffD2)
apply (rule proj_in_iff)
apply (rule equiv_lsbis2[OF coalg_T])
apply assumption

```

```

apply (rule conjI)
apply (rule ballI)
apply (rule sym)
apply (rule trans)
apply (rule univ_commute)
apply (rule equiv_lsbis1[OF coalg_T])
apply (rule congruent_strQ1)
apply assumption
apply (rule o_apply)

```

```

apply (rule ballI)
apply (rule sym)
apply (rule trans)
apply (rule univ_commute)
apply (rule equiv_lsbis2[OF coalg_T])
apply (rule congruent_strQ2)
apply assumption
apply (rule o_apply)
done

```

```

lemmas mor_final = mor_comp[OF mor_beh mor_T_final]
lemmas in_car_final1 = mor_image1'[OF mor_final UNIV_I]
lemmas in_car_final2 = mor_image2'[OF mor_final UNIV_I]

```

```

typedef (overloaded) 'a JF1 = car_final1 :: 'a carrier set set
by (rule exI) (rule in_car_final1)

```

```

typedef (overloaded) 'a JF2 = car_final2 :: 'a carrier set set
by (rule exI) (rule in_car_final2)

```

definition *dtor1* **where**

```

dtor1 x = F1map id Abs_JF1 Abs_JF2 (str_final1 (Rep_JF1 x))

```

definition *dtor2* **where**

```

dtor2 x = F2map id Abs_JF1 Abs_JF2 (str_final2 (Rep_JF2 x))

```

lemma *mor_Rep_JF*: *mor UNIV UNIV dtor1 dtor2*
car_final1 car_final2 str_final1 str_final2
Rep_JF1 Rep_JF2

unfolding *mor_def dtor1_def dtor2_def*
apply (*rule conjI*)
apply (*rule conjI*)
apply (*rule ballI*)
apply (*rule Rep_JF1*)
apply (*rule ballI*)
apply (*rule Rep_JF2*)

apply (*rule conjI*)
apply (*rule ballI*)
apply (*rule trans[OF F1map_comp_id]*)
apply (*rule F1map_congL*)
apply (*rule ballI*)
apply (*rule trans[OF o_apply]*)
apply (*rule Abs_JF1_inverse*)
apply (*erule set_rev_mp*)
apply (*rule coalg_F1set2*)
apply (*rule coalg_final*)
apply (*rule Rep_JF1*)
apply (*rule ballI*)
apply (*rule trans[OF o_apply]*)
apply (*rule Abs_JF2_inverse*)
apply (*erule set_rev_mp*)
apply (*rule coalg_F1set3*)
apply (*rule coalg_final*)
apply (*rule Rep_JF1*)

apply (*rule ballI*)
apply (*rule trans[OF F2map_comp_id]*)
apply (*rule F2map_congL*)
apply (*rule ballI*)
apply (*rule trans[OF o_apply]*)
apply (*rule Abs_JF1_inverse*)
apply (*erule set_rev_mp*)
apply (*rule coalg_F2set2*)
apply (*rule coalg_final*)
apply (*rule Rep_JF2*)
apply (*rule ballI*)
apply (*rule trans[OF o_apply]*)
apply (*rule Abs_JF2_inverse*)
apply (*erule set_rev_mp*)
apply (*rule coalg_F2set3*)
apply (*rule coalg_final*)
apply (*rule Rep_JF2*)
done

lemma *mor_Abs_JF*: *mor car_final1 car_final2 str_final1 str_final2*
UNIV UNIV dtor1 dtor2 Abs_JF1 Abs_JF2

unfolding *mor_def dtor1_def dtor2_def*
apply (*rule conjI*)
apply (*rule conjI*)
apply (*rule ballI*)
apply (*rule UNIV_I*)
apply (*rule ballI*)
apply (*rule UNIV_I*)

apply (*rule conjI*)
apply (*rule ballI*)
apply (*erule sym[OF arg_cong[OF Abs_JF1_inverse]]*)

```

apply (rule ballI)
apply (erule sym[OF arg_cong[OF Abs_JF2_inverse]])
done

```

definition *unfold1* **where**

```

unfold1 s1 s2 x =
  Abs_JF1 ((Equiv_Relations.proj (lsbis1 carT1 carT2 strT1 strT2) o beh1 s1 s2) x)

```

definition *unfold2* **where**

```

unfold2 s1 s2 x =
  Abs_JF2 ((Equiv_Relations.proj (lsbis2 carT1 carT2 strT1 strT2) o beh2 s1 s2) x)

```

lemma *mor_unfold*:

```

mor UNIV UNIV s1 s2 UNIV UNIV dtor1 dtor2 (unfold1 s1 s2) (unfold2 s1 s2)

```

```

apply (rule iffD2)
apply (rule mor_UNIV)
apply (rule conjI)
apply (rule ext)
apply (rule sym[OF trans[OF o_apply]])
apply (rule trans[OF dtor1_def])
apply (rule trans[OF arg_cong[OF unfold1_def]])
apply (rule trans[OF arg_cong[OF Abs_JF1_inverse[OF in_car_final1]])]
apply (rule trans[OF arg_cong[OF sym[OF morE1[OF mor_final UNIV_I]])]]
apply (rule trans[OF F1map_comp_id])
apply (rule sym[OF trans[OF o_apply]])
apply (rule F1.map_cong0)
apply (rule refl)
apply (rule trans[OF unfold1_def])
apply (rule sym[OF o_apply])
apply (rule trans[OF unfold2_def])
apply (rule sym[OF o_apply])

```

```

apply (rule ext)
apply (rule sym[OF trans[OF o_apply]])
apply (rule trans[OF dtor2_def])
apply (rule trans[OF arg_cong[OF unfold2_def]])
apply (rule trans[OF arg_cong[OF Abs_JF2_inverse[OF in_car_final2]])]
apply (rule trans[OF arg_cong[OF sym[OF morE2[OF mor_final UNIV_I]])]]
apply (rule trans[OF F2map_comp_id])
apply (rule sym[OF trans[OF o_apply]])
apply (rule F2.map_cong0)
apply (rule refl)
apply (rule trans[OF unfold1_def])
apply (rule sym[OF o_apply])
apply (rule trans[OF unfold2_def])
apply (rule sym[OF o_apply])
done

```

lemmas *unfold1* = sym[OF morE1[OF mor_unfold UNIV_I]]

lemmas *unfold2* = sym[OF morE2[OF mor_unfold UNIV_I]]

lemma *JF_cind*: $sbis\ UNIV\ UNIV\ dtor1\ dtor2\ R1\ R2 \implies R1 \subseteq Id \wedge R2 \subseteq Id$

```

apply (rule rev_mp)
apply (tactic << forward_tac @ {context} @ {thms bis_def[THEN iffD1]} 1 >>)
apply (erule conjE)+
apply (rule bis_cong)
apply (rule bis_Comp)
apply (rule bis_converse)
apply (rule bis_Gr)
apply (rule tcoalg)
apply (rule mor_Rep_JF)
apply (rule bis_Comp)
apply assumption
apply (rule bis_Gr)

```

```

    apply (rule tcoalg)
    apply (rule mor_Rep_JF)
    apply (erule relImage_Gr)
    apply (erule relImage_Gr)

apply (rule impI)
apply (rule rev_mp)
apply (rule bis_cong)
  apply (rule bis_Comp)
  apply (rule bis_Gr)
  apply (rule coalg_T)
  apply (rule mor_T_final)
  apply (rule bis_Comp)
  apply (rule sbis_lsbis)
  apply (rule bis_converse)
  apply (rule bis_Gr)
  apply (rule coalg_T)
  apply (rule mor_T_final)
  apply (rule relInvImage_Gr[OF lsbis1_incl])
  apply (rule relInvImage_Gr[OF lsbis2_incl])

```

```

apply (rule impI)
apply (rule conjI)
  apply (rule subset_trans)
  apply (rule relInvImage_UNIV_relImage)
  apply (rule subset_trans)
  apply (rule relInvImage_mono)
  apply (rule subset_trans)
  apply (erule incl_lsbis1)
  apply (rule ord_eq_le_trans)
  apply (rule sym[OF relImage_relInvImage])
  apply (rule xt1(3))
  apply (rule Sigma_cong)
  apply (rule proj_image)
  apply (rule proj_image)
  apply (rule lsbis1_incl)
  apply (rule subset_trans)
  apply (rule relImage_mono)
  apply (rule incl_lsbis1)
  apply assumption
  apply (rule relImage_proj)
  apply (rule equiv_lsbis1[OF coalg_T])
  apply (rule relInvImage_Id_on)
  apply (rule Rep_JF1_inject)

```

```

apply (rule subset_trans)
  apply (rule relInvImage_UNIV_relImage)
  apply (rule subset_trans)
  apply (rule relInvImage_mono)
  apply (rule subset_trans)
  apply (erule incl_lsbis2)
  apply (rule ord_eq_le_trans)
  apply (rule sym[OF relImage_relInvImage])
  apply (rule xt1(3))
  apply (rule Sigma_cong)
  apply (rule proj_image)
  apply (rule proj_image)
  apply (rule lsbis2_incl)
  apply (rule subset_trans)
  apply (rule relImage_mono)
  apply (rule incl_lsbis2)
  apply assumption
  apply (rule relImage_proj)

```



```

  apply (rule equiv_lsbis2[OF coalg_T])
  apply (rule relInvImage_Id_on)
  apply (rule Rep_JF2_inject)
done

lemmas JF_cind1 = conjunct1[OF JF_cind]
lemmas JF_cind2 = conjunct2[OF JF_cind]

lemma unfold_unique_mor:
  mor UNIV UNIV s1 s2 UNIV UNIV dtor1 dtor2 f1 f2  $\implies$ 
  f1 = unfold1 s1 s2  $\wedge$  f2 = unfold2 s1 s2
  apply (rule conjI)
  apply (rule ext)
  apply (erule IdD[OF set_mp[OF JF_cind1[OF bis_image2[OF tcoalg _ tcoalg]]]])
  apply (rule mor_comp[OF mor_final mor_Abs_JF])
  apply (rule image2_eqI)
  apply (rule refl)
  apply (rule trans[OF arg_cong[OF unfold1_def]])
  apply (rule sym[OF o_apply])
  apply (rule UNIV_I)

  apply (rule ext)
  apply (erule IdD[OF set_mp[OF JF_cind2[OF bis_image2[OF tcoalg _ tcoalg]]]])
  apply (rule mor_comp[OF mor_final mor_Abs_JF])
  apply (rule image2_eqI)
  apply (rule refl)
  apply (rule trans[OF arg_cong[OF unfold2_def]])
  apply (rule sym[OF o_apply])
  apply (rule UNIV_I)
done

lemmas unfold_unique = unfold_unique_mor[OF iffD2[OF mor_UNIV], OF conjI]
lemmas unfold1_dtor = sym[OF conjunct1[OF unfold_unique_mor[OF mor_id]]]
lemmas unfold2_dtor = sym[OF conjunct2[OF unfold_unique_mor[OF mor_id]]]

lemmas unfold1_o_dtor1 =
  trans[OF conjunct1[OF unfold_unique_mor[OF mor_comp[OF mor_str mor_unfold]]] unfold1_dtor]
lemmas unfold2_o_dtor2 =
  trans[OF conjunct2[OF unfold_unique_mor[OF mor_comp[OF mor_str mor_unfold]]] unfold2_dtor]

definition ctor1 where ctor1 = unfold1 (F1map id dtor1 dtor2) (F2map id dtor1 dtor2)
definition ctor2 where ctor2 = unfold2 (F1map id dtor1 dtor2) (F2map id dtor1 dtor2)

lemma ctor1_o_dtor1:
  ctor1 o dtor1 = id
  unfolding ctor1_def
  apply (rule unfold1_o_dtor1)
done

lemma ctor2_o_dtor2:
  ctor2 o dtor2 = id
  unfolding ctor2_def
  apply (rule unfold2_o_dtor2)
done

lemma dtor1_o_ctor1:
  dtor1 o ctor1 = id
  unfolding ctor1_def
  apply (rule ext)
  apply (rule trans[OF o_apply])
  apply (rule trans[OF unfold1])
  apply (rule trans[OF F1map_comp_id])

```

```

apply (rule trans[OF F1map_congL])
  apply (rule ballI)
  apply (rule trans[OF fun_cong[OF unfold1_o_dtor1] id_apply])
  apply (rule ballI)
  apply (rule trans[OF fun_cong[OF unfold2_o_dtor2] id_apply])
apply (rule sym[OF id_apply])
done

```

```

lemma dtor2_o_ctor2:
  dtor2 o ctor2 = id
  unfolding ctor2_def
  apply (rule ext)
  apply (rule trans[OF o_apply])
  apply (rule trans[OF unfold2])
  apply (rule trans[OF F2map_comp_id])
  apply (rule trans[OF F2map_congL])
  apply (rule ballI)
  apply (rule trans[OF fun_cong[OF unfold1_o_dtor1] id_apply])
  apply (rule ballI)
  apply (rule trans[OF fun_cong[OF unfold2_o_dtor2] id_apply])
  apply (rule sym[OF id_apply])
done

```

```

lemmas dtor1_ctor1 = pointfree_idE[OF dtor1_o_ctor1]
lemmas dtor2_ctor2 = pointfree_idE[OF dtor2_o_ctor2]
lemmas ctor1_dtor1 = pointfree_idE[OF ctor1_o_dtor1]
lemmas ctor2_dtor2 = pointfree_idE[OF ctor2_o_dtor2]

```

```

lemmas bij_dtor1 = o_bij[OF ctor1_o_dtor1 dtor1_o_ctor1]
lemmas inj_dtor1 = bij_is_inj[OF bij_dtor1]
lemmas surj_dtor1 = bij_is_surj[OF bij_dtor1]
lemmas dtor1_nchotomy = surjD[OF surj_dtor1]
lemmas dtor1_diff = inj_eq[OF inj_dtor1]
lemmas dtor1_cases = exE[OF dtor1_nchotomy]
lemmas bij_dtor2 = o_bij[OF ctor2_o_dtor2 dtor2_o_ctor2]
lemmas inj_dtor2 = bij_is_inj[OF bij_dtor2]
lemmas surj_dtor2 = bij_is_surj[OF bij_dtor2]
lemmas dtor2_nchotomy = surjD[OF surj_dtor2]
lemmas dtor2_diff = inj_eq[OF inj_dtor2]
lemmas dtor2_cases = exE[OF dtor2_nchotomy]

```

```

lemmas bij_ctor1 = o_bij[OF dtor1_o_ctor1 ctor1_o_dtor1]
lemmas inj_ctor1 = bij_is_inj[OF bij_ctor1]
lemmas surj_ctor1 = bij_is_surj[OF bij_ctor1]
lemmas ctor1_nchotomy = surjD[OF surj_ctor1]
lemmas ctor1_diff = inj_eq[OF inj_ctor1]
lemmas ctor1_cases = exE[OF ctor1_nchotomy]
lemmas bij_ctor2 = o_bij[OF dtor2_o_ctor2 ctor2_o_dtor2]
lemmas inj_ctor2 = bij_is_inj[OF bij_ctor2]
lemmas surj_ctor2 = bij_is_surj[OF bij_ctor2]
lemmas ctor2_nchotomy = surjD[OF surj_ctor2]
lemmas ctor2_diff = inj_eq[OF inj_ctor2]
lemmas ctor2_cases = exE[OF ctor2_nchotomy]

```

```

lemmas ctor1_unfold1 = iffD1[OF dtor1_diff trans[OF unfold1 sym[OF dtor1_ctor1]]]
lemmas ctor2_unfold2 = iffD1[OF dtor2_diff trans[OF unfold2 sym[OF dtor2_ctor2]]]

```

```

definition corec1 where corec1 s1 s2 =
  unfold1 (case_sum (F1map id Inl Inl o dtor1) s1)
    (case_sum (F2map id Inl Inl o dtor2) s2) o Inr
definition corec2 where corec2 s1 s2 =

```

unfold2 (case_sum (F1map id Inl Inl o dtor1) s1)
(case_sum (F2map id Inl Inl o dtor2) s2) o Inr

lemma dtor1_o_unfold1: dtor1 o unfold1 s1 s2 = F1map id (unfold1 s1 s2) (unfold2 s1 s2) o s1
by (tactic ⟨rtac @{context} (BNF_Tactics.mk_pointfree2 @{context} @{thm unfold1}) 1⟩)

lemma dtor2_o_unfold2: dtor2 o unfold2 s1 s2 = F2map id (unfold1 s1 s2) (unfold2 s1 s2) o s2
by (tactic ⟨rtac @{context} (BNF_Tactics.mk_pointfree2 @{context} @{thm unfold2}) 1⟩)

lemma corec1_Inl_sum:

unfold1 (case_sum (F1map id Inl Inl o dtor1) s1) (case_sum (F2map id Inl Inl o dtor2) s2) o Inl = id
apply (rule trans[OF conjunct1[OF unfold_unique] unfold1_dtor])
apply (rule trans[OF arg_cong2[of _ _ _ _ (o), OF F1.map_comp0[of id, unfolded id_o] refl]])
apply (rule sym[OF trans[OF o_assoc]])
apply (rule trans[OF arg_cong2[of _ _ _ _ (o), OF dtor1_o_unfold1 refl]])
apply (rule box_equals[OF _ o_assoc o_assoc])
apply (rule arg_cong2[of _ _ _ _ (o), OF refl case_sum_o_inj(1)])
apply (rule trans[OF arg_cong2[of _ _ _ _ (o), OF F2.map_comp0[of id, unfolded id_o] refl]])
apply (rule sym[OF trans[OF o_assoc]])
apply (rule trans[OF arg_cong2[of _ _ _ _ (o), OF dtor2_o_unfold2 refl]])
apply (rule box_equals[OF _ o_assoc o_assoc])
apply (rule arg_cong2[of _ _ _ _ (o), OF refl case_sum_o_inj(1)])
done

lemma corec2_Inl_sum:

unfold2 (case_sum (F1map id Inl Inl o dtor1) s1) (case_sum (F2map id Inl Inl o dtor2) s2) o Inl = id
apply (rule trans[OF conjunct2[OF unfold_unique] unfold2_dtor])
apply (rule trans[OF arg_cong2[of _ _ _ _ (o), OF F1.map_comp0[of id, unfolded id_o] refl]])
apply (rule sym[OF trans[OF o_assoc]])
apply (rule trans[OF arg_cong2[of _ _ _ _ (o), OF dtor1_o_unfold1 refl]])
apply (rule box_equals[OF _ o_assoc o_assoc])
apply (rule arg_cong2[of _ _ _ _ (o), OF refl case_sum_o_inj(1)])
apply (rule trans[OF arg_cong2[of _ _ _ _ (o), OF F2.map_comp0[of id, unfolded id_o] refl]])
apply (rule sym[OF trans[OF o_assoc]])
apply (rule trans[OF arg_cong2[of _ _ _ _ (o), OF dtor2_o_unfold2 refl]])
apply (rule box_equals[OF _ o_assoc o_assoc])
apply (rule arg_cong2[of _ _ _ _ (o), OF refl case_sum_o_inj(1)])
done

lemma case_sum_expand_Inr: f o Inl = g \implies case_sum g (f o Inr) = f
by (auto split: sum.splits)

theorem corec1:

dtor1 (corec1 s1 s2 a) =
F1map id (case_sum id (corec1 s1 s2)) (case_sum id (corec2 s1 s2)) (s1 a)
unfolding corec1_def corec2_def o_apply unfold1 sum.case
case_sum_expand_Inr[OF corec1_Inl_sum] case_sum_expand_Inr[OF corec2_Inl_sum] ..

theorem corec2:

dtor2 (corec2 s1 s2 a) =
F2map id (case_sum id (corec1 s1 s2)) (case_sum id (corec2 s1 s2)) (s2 a)
unfolding corec1_def corec2_def o_apply unfold2 sum.case
case_sum_expand_Inr[OF corec1_Inl_sum] case_sum_expand_Inr[OF corec2_Inl_sum] ..

lemma corec_unique:

F1map id (case_sum id f1) (case_sum id f2) o s1 = dtor1 o f1 \implies
F2map id (case_sum id f1) (case_sum id f2) o s2 = dtor2 o f2 \implies
f1 = corec1 s1 s2 \wedge f2 = corec2 s1 s2

unfolding corec1_def corec2_def case_sum_expand_Inr'[OF corec1_Inl_sum] case_sum_expand_Inr'[OF corec2_Inl_sum]
apply (rule unfold_unique)
apply (unfold o_case_sum_id_o_id F1.map_comp0[symmetric] F2.map_comp0[symmetric]
F1.map_id0 F2.map_id0 o_assoc case_sum_o_inj(1))
apply (erule arg_cong2[of _ _ _ _ case_sum, OF refl])
apply (erule arg_cong2[of _ _ _ _ case_sum, OF refl])

done

2.7 Coinduction

lemma *Frel_coind*:

$$\llbracket \forall a b. \text{phi1 } a b \longrightarrow F1\text{rel } (=) \text{ phi1 phi2 } (\text{dtor1 } a) (\text{dtor1 } b);$$
$$\forall a b. \text{phi2 } a b \longrightarrow F2\text{rel } (=) \text{ phi1 phi2 } (\text{dtor2 } a) (\text{dtor2 } b) \rrbracket \implies$$
$$(\text{phi1 } a1 b1 \longrightarrow a1 = b1) \wedge (\text{phi2 } a2 b2 \longrightarrow a2 = b2)$$

apply (*rule rev_mp*)

apply (*rule JF_cind*)

apply (*rule iffD2*)

apply (*rule bis_Frel*)

apply (*rule conjI*)

apply (*rule conjI*)

apply (*rule ord_le_eq_trans*[*OF subset_UNIV UNIV_Times_UNIV*[*THEN sym*]])

apply (*rule ord_le_eq_trans*[*OF subset_UNIV UNIV_Times_UNIV*[*THEN sym*]])

apply (*rule conjI*)

apply (*rule allI*)⁺

apply (*rule impI*)

apply (*erule allE*)⁺

apply (*rule predicate2D*[*OF eq_refl*[*OF F1rel_cong*]])

apply (*rule refl*)

apply (*rule in_rel_Collect_case_prod_eq*[*symmetric*])

apply (*rule in_rel_Collect_case_prod_eq*[*symmetric*])

apply (*erule mp*)

apply (*erule CollectE*)

apply (*erule case_prodD*)

apply (*rule allI*)⁺

apply (*rule impI*)

apply (*erule allE*)⁺

apply (*rule predicate2D*[*OF eq_refl*[*OF F2rel_cong*]])

apply (*rule refl*)

apply (*rule in_rel_Collect_case_prod_eq*[*symmetric*])

apply (*rule in_rel_Collect_case_prod_eq*[*symmetric*])

apply (*erule mp*)

apply (*erule CollectE*)

apply (*erule case_prodD*)

apply (*rule impI*)

apply (*erule conjE*)⁺

apply (*rule conjI*)

apply (*rule impI*)

apply (*rule IdD*)

apply (*erule set_mp*)

apply (*rule CollectI*)

apply (*erule case_prodI*)

apply (*rule impI*)

apply (*rule IdD*)

apply (*erule set_mp*)

apply (*rule CollectI*)

apply (*erule case_prodI*)

done

2.8 The Result as an BNF

abbreviation *JF1map* **where**

$$JF1map \ u \equiv \text{unfold1 } (F1map \ u \ \text{id} \ \text{id} \ o \ \text{dtor1}) \ (F2map \ u \ \text{id} \ \text{id} \ o \ \text{dtor2})$$

abbreviation *JF2map* **where**

$JF2map\ u \equiv unfold2\ (F1map\ u\ id\ id\ o\ dtor1)\ (F2map\ u\ id\ id\ o\ dtor2)$

lemma *JF1map*: $dtor1\ o\ JF1map\ u = F1map\ u\ (JF1map\ u)\ (JF2map\ u)\ o\ dtor1$

apply (rule *ext*)
apply (rule *sym*[*OF trans*[*OF o_apply*]])
apply (rule *sym*[*OF trans*[*OF o_apply*]])
apply (rule *trans*[*OF unfold1*])
apply (rule *box_equals*[*OF F1.map_comp _ F1.map_cong0, rotated*])
apply (rule *fun_cong*[*OF id_o*])
apply (rule *fun_cong*[*OF o_id*])
apply (rule *fun_cong*[*OF o_id*])
apply (rule *sym*[*OF arg_cong*[*OF o_apply*]])
done

lemma *JF2map*: $dtor2\ o\ JF2map\ u = F2map\ u\ (JF1map\ u)\ (JF2map\ u)\ o\ dtor2$

apply (rule *ext*)
apply (rule *sym*[*OF trans*[*OF o_apply*]])
apply (rule *sym*[*OF trans*[*OF o_apply*]])
apply (rule *trans*[*OF unfold2*])
apply (rule *box_equals*[*OF F2.map_comp _ F2.map_cong0, rotated*])
apply (rule *fun_cong*[*OF id_o*])
apply (rule *fun_cong*[*OF o_id*])
apply (rule *fun_cong*[*OF o_id*])
apply (rule *sym*[*OF arg_cong*[*OF o_apply*]])
done

lemmas *JF1map_simps* = *o_eq_dest*[*OF JF1map*]

lemmas *JF2map_simps* = *o_eq_dest*[*OF JF2map*]

theorem *JF1map_id*: $JF1map\ id = id$

apply (rule *trans*)
apply (rule *conjunct1*)
apply (rule *unfold_unique*)
apply (rule *sym*[*OF JF1map*])
apply (rule *sym*[*OF JF2map*])
apply (rule *unfold1_dtor*)
done

theorem *JF2map_id*: $JF2map\ id = id$

apply (rule *trans*)
apply (rule *conjunct2*)
apply (rule *unfold_unique*)
apply (rule *sym*[*OF JF1map*])
apply (rule *sym*[*OF JF2map*])
apply (rule *unfold2_dtor*)
done

lemma *JFmap_unique*:

$[[dtor1\ o\ u = F1map\ f\ u\ v\ o\ dtor1; dtor2\ o\ v = F2map\ f\ u\ v\ o\ dtor2]] \implies$

$u = JF1map\ f \wedge v = JF2map\ f$

apply (rule *unfold_unique*)
unfolding *o_assoc* *F1.map_comp0[symmetric]* *F2.map_comp0[symmetric]* *id_o_o_id*
apply (rule *sym*)
apply (rule *sym*)
done

theorem *JF1map_comp*: $JF1map\ (g\ o\ f) = JF1map\ g\ o\ JF1map\ f$

apply (rule *sym*)
apply (rule *conjunct1*)
apply (rule *JFmap_unique*)
apply (rule *trans*[*OF o_assoc*])
apply (rule *trans*[*OF arg_cong2*[*of _ _ _ (o), OF JF1map refl*]])
apply (rule *trans*[*OF sym*[*OF o_assoc*]])

```

apply (rule trans[OF arg_cong[OF JF1map]])
apply (rule trans[OF o_assoc])
apply (rule arg_cong2[of _ _ _ (o), OF sym[OF F1.map_comp0] refl])

```

```

apply (rule trans[OF o_assoc])
apply (rule trans[OF arg_cong2[of _ _ _ (o), OF JF2map refl]])
apply (rule trans[OF sym[OF o_assoc]])
apply (rule trans[OF arg_cong[OF JF2map]])
apply (rule trans[OF o_assoc])
apply (rule arg_cong2[of _ _ _ (o), OF sym[OF F2.map_comp0] refl])
done

```

theorem *JF2map_comp*: $JF2map (g \circ f) = JF2map g \circ JF2map f$

```

apply (rule sym)
apply (rule conjunct2)
apply (tactic << rtac @{context} (Thm.permute_prem0 1 @{thm unfold_unique}) 1 >>)

```

```

apply (rule trans[OF o_assoc])
apply (rule trans[OF arg_cong[OF sym[OF F2.map_comp0]]])
apply (rule sym[OF trans[OF o_assoc]])
apply (rule trans[OF arg_cong2[OF JF2map refl]])
apply (rule trans[OF sym[OF o_assoc]])
apply (rule trans[OF arg_cong[OF JF2map]])
apply (rule trans[OF o_assoc])
apply (rule trans[OF arg_cong2[OF sym[OF F2.map_comp0] refl]])
apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule sym)
apply (rule trans[OF o_apply])
apply (rule F2.map_comp0)
  apply (rule trans[OF o_apply])
  apply (rule id_apply)
apply (rule trans[OF o_apply])
apply (rule arg_cong[OF id_apply])
apply (rule trans[OF o_apply])
apply (rule arg_cong[OF id_apply])

```

```

apply (rule trans[OF o_assoc])
apply (rule trans[OF arg_cong[OF sym[OF F1.map_comp0]]])
apply (rule sym[OF trans[OF o_assoc]])
apply (rule trans[OF arg_cong2[OF JF1map refl]])
apply (rule trans[OF sym[OF o_assoc]])
apply (rule trans[OF arg_cong[OF JF1map]])
apply (rule trans[OF o_assoc])
apply (rule trans[OF arg_cong2[OF sym[OF F1.map_comp0] refl]])
apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule sym)
apply (rule trans[OF o_apply])
apply (rule F1.map_comp0)
  apply (rule trans[OF o_apply])
  apply (rule id_apply)
apply (rule trans[OF o_apply])
apply (rule arg_cong[OF id_apply])
apply (rule trans[OF o_apply])
apply (rule arg_cong[OF id_apply])
done

```

definition *JFcol* **where**

```

JFcol = rec_nat (%a. {}, %b. {})
  (%n rec.
  (%a. F1set1 (dctor1 a)  $\cup$ 

```

$$\begin{aligned}
& ((\bigcup a' \in F1set2 \ (dtor1 \ a). \ fst \ rec \ a') \cup \\
& (\bigcup a' \in F1set3 \ (dtor1 \ a). \ snd \ rec \ a')), \\
\%b. & F2set1 \ (dtor2 \ b) \cup \\
& ((\bigcup b' \in F2set2 \ (dtor2 \ b). \ fst \ rec \ b') \cup \\
& (\bigcup b' \in F2set3 \ (dtor2 \ b). \ snd \ rec \ b')))
\end{aligned}$$

abbreviation *JF1col* **where** $JF1col \ n \equiv \text{fst} \ (JFcol \ n)$

abbreviation *JF2col* **where** $JF2col \ n \equiv \text{snd} \ (JFcol \ n)$

lemmas $JF1col_0 = \text{fun_cong}[OF \ \text{fstI}[OF \ \text{rec_nat_0_imp}[OF \ JFcol_def]]]$

lemmas $JF2col_0 = \text{fun_cong}[OF \ \text{sndI}[OF \ \text{rec_nat_0_imp}[OF \ JFcol_def]]]$

lemmas $JF1col_Suc = \text{fun_cong}[OF \ \text{fstI}[OF \ \text{rec_nat_Suc_imp}[OF \ JFcol_def]]]$

lemmas $JF2col_Suc = \text{fun_cong}[OF \ \text{sndI}[OF \ \text{rec_nat_Suc_imp}[OF \ JFcol_def]]]$

lemma *JFcol_minimal*:

$$\begin{aligned}
& \llbracket \bigwedge a. \ F1set1 \ (dtor1 \ a) \subseteq K1 \ a; \\
& \bigwedge b. \ F2set1 \ (dtor2 \ b) \subseteq K2 \ b; \\
& \bigwedge a \ a'. \ a' \in F1set2 \ (dtor1 \ a) \implies K1 \ a' \subseteq K1 \ a; \\
& \bigwedge a \ b'. \ b' \in F1set3 \ (dtor1 \ a) \implies K2 \ b' \subseteq K1 \ a; \\
& \bigwedge b \ a'. \ a' \in F2set2 \ (dtor2 \ b) \implies K1 \ a' \subseteq K2 \ b; \\
& \bigwedge b \ b'. \ b' \in F2set3 \ (dtor2 \ b) \implies K2 \ b' \subseteq K2 \ b \rrbracket \implies
\end{aligned}$$

$\forall a \ b. \ JF1col \ n \ a \subseteq K1 \ a \wedge JF2col \ n \ b \subseteq K2 \ b$

apply (*rule nat_induct*)

apply (*rule allI*)⁺

apply (*rule conjI*)

apply (*rule ord_eq_le_trans*)

apply (*rule JF1col_0*)

apply (*rule empty_subsetI*)

apply (*rule ord_eq_le_trans*)

apply (*rule JF2col_0*)

apply (*rule empty_subsetI*)

apply (*rule allI*)⁺

apply (*rule conjI*)

apply (*rule ord_eq_le_trans*)

apply (*rule JF1col_Suc*)

apply (*rule Un_least*)

apply *assumption*

apply (*rule Un_least*)

apply (*rule UN_least*)

apply (*erule allE conjE*)⁺

apply (*rule subset_trans*)

apply *assumption*

apply *assumption*

apply (*rule UN_least*)

apply (*erule allE conjE*)⁺

apply (*rule subset_trans*)

apply *assumption*

apply *assumption*

apply (*rule ord_eq_le_trans*)

apply (*rule JF2col_Suc*)

apply (*rule Un_least*)

apply *assumption*

apply (*rule Un_least*)

apply (*rule UN_least*)

apply (*erule allE conjE*)⁺

apply (*rule subset_trans*)

apply *assumption*

apply *assumption*

apply (rule *UN_least*)
apply (erule *allE conjE*)
apply (rule *subset_trans*)
apply *assumption*
apply *assumption*
done

lemma *JFset_minimal*:

$\llbracket \bigwedge a. F1set1 (dtor1 a) \subseteq K1 a;$
 $\bigwedge b. F2set1 (dtor2 b) \subseteq K2 b;$
 $\bigwedge a a'. a' \in F1set2 (dtor1 a) \implies K1 a' \subseteq K1 a;$
 $\bigwedge a b'. b' \in F1set3 (dtor1 a) \implies K2 b' \subseteq K1 a;$
 $\bigwedge b a'. a' \in F2set2 (dtor2 b) \implies K1 a' \subseteq K2 b;$
 $\bigwedge b b'. b' \in F2set3 (dtor2 b) \implies K2 b' \subseteq K2 b \rrbracket \implies$
 $(\bigcup n. JF1col n a) \subseteq K1 a \wedge (\bigcup n. JF2col n b) \subseteq K2 b$
apply (rule *conjI*)
apply (rule *UN_least*)
apply (rule *rev_mp*)
apply (rule *JFcol_minimal*)
apply *assumption*
apply *assumption*
apply *assumption*
apply *assumption*
apply *assumption*
apply *assumption*
apply (rule *impI*)
apply (erule *allE conjE*)
apply *assumption*

apply (rule *UN_least*)
apply (rule *rev_mp*)
apply (rule *JFcol_minimal*)
apply *assumption*
apply *assumption*
apply *assumption*
apply *assumption*
apply *assumption*
apply (rule *impI*)
apply (erule *allE conjE*)
apply *assumption*
done

abbreviation *JF1set* **where** $JF1set a \equiv (\bigcup n. JF1col n a)$

abbreviation *JF2set* **where** $JF2set a \equiv (\bigcup n. JF2col n a)$

lemma *F1set1_incl_JF1set*:

$F1set1 (dtor1 a) \subseteq JF1set a$
apply (rule *SUP_upper2*)
apply (rule *UNIV_I*)
apply (rule *ord_le_eq_trans*)
apply (rule *Un_upper1*)
apply (rule *sym*)
apply (rule *JF1col_Suc*)
done

lemma *F2set1_incl_JF2set*:

$F2set1 (dtor2 a) \subseteq JF2set a$
apply (rule *SUP_upper2*)
apply (rule *UNIV_I*)
apply (rule *ord_le_eq_trans*)


```

  apply (rule Un_upper1)
  apply (rule sym)
  apply (rule JF2col_Suc)
done

```

```

lemma F1set2_JF1set_incl_JF1set:
  a' ∈ F1set2 (dtor1 a) ⇒ JF1set a' ⊆ JF1set a
  apply (rule UN_least)
  apply (rule subsetI)
  apply (rule UN_I)
  apply (rule UNIV_I)
  apply (rule set_mp)
  apply (rule equalityD2)
  apply (rule JF1col_Suc)
  apply (rule UnI2)
  apply (tactic ⟨⟨ rtac @{context} (BNF_Util.mk_UnIN 2 1) 1 ⟩⟩)
  apply (erule UN_I)
  apply assumption
done

```

```

lemma F1set3_JF2set_incl_JF1set:
  a' ∈ F1set3 (dtor1 a) ⇒ JF2set a' ⊆ JF1set a
  apply (rule UN_least)
  apply (rule subsetI)
  apply (rule UN_I)
  apply (rule UNIV_I)
  apply (rule set_mp)
  apply (rule equalityD2)
  apply (rule JF1col_Suc)
  apply (rule UnI2)
  apply (tactic ⟨⟨ rtac @{context} (BNF_Util.mk_UnIN 2 2) 1 ⟩⟩)
  apply (erule UN_I)
  apply assumption
done

```

```

lemma F2set2_JF1set_incl_JF2set:
  a' ∈ F2set2 (dtor2 a) ⇒ JF1set a' ⊆ JF2set a
  apply (rule UN_least)
  apply (rule subsetI)
  apply (rule UN_I)
  apply (rule UNIV_I)
  apply (rule set_mp)
  apply (rule equalityD2)
  apply (rule JF2col_Suc)
  apply (rule UnI2)
  apply (tactic ⟨⟨ rtac @{context} (BNF_Util.mk_UnIN 2 1) 1 ⟩⟩)
  apply (erule UN_I)
  apply assumption
done

```

```

lemma F2set3_JF2set_incl_JF2set:
  a' ∈ F2set3 (dtor2 a) ⇒ JF2set a' ⊆ JF2set a
  apply (rule UN_least)
  apply (rule subsetI)
  apply (rule UN_I)
  apply (rule UNIV_I)
  apply (rule set_mp)
  apply (rule equalityD2)
  apply (rule JF2col_Suc)
  apply (rule UnI2)
  apply (tactic ⟨⟨ rtac @{context} (BNF_Util.mk_UnIN 2 2) 1 ⟩⟩)
  apply (erule UN_I)
  apply assumption

```

done

lemmas $F1set1_JF1set = set_mp[OF F1set1_incl_JF1set]$
lemmas $F2set1_JF2set = set_mp[OF F2set1_incl_JF2set]$
lemmas $F1set2_JF1set_JF1set = set_mp[OF F1set2_JF1set_incl_JF1set]$
lemmas $F1set3_JF2set_JF1set = set_mp[OF F1set3_JF2set_incl_JF1set]$
lemmas $F2set2_JF1set_JF2set = set_mp[OF F2set2_JF1set_incl_JF2set]$
lemmas $F2set3_JF2set_JF2set = set_mp[OF F2set3_JF2set_incl_JF2set]$

lemma $JFset_le$:

fixes $a :: 'a JF1$ and $b :: 'a JF2$

shows

$JF1set\ a \subseteq F1set1\ (dtor1\ a) \cup (UNION\ (F1set2\ (dtor1\ a))\ JF1set \cup UNION\ (F1set3\ (dtor1\ a))\ JF2set) \wedge$
 $JF2set\ b \subseteq F2set1\ (dtor2\ b) \cup (UNION\ (F2set2\ (dtor2\ b))\ JF1set \cup UNION\ (F2set3\ (dtor2\ b))\ JF2set)$

apply (rule $JFset_minimal$)
 apply (rule $\bar{U}n_upper1$)
 apply (rule Un_upper1)
 apply (rule $subsetI$)
 apply (rule $UnI2$)
 apply (tactic $\langle\langle rtac\ @\{context\}\ (BNF_Util.mk_UnIN\ 2\ 1)\ 1\ \rangle\rangle$)
 apply (erule UN_I)
 apply (erule UnE)
 apply (erule $F1set1_JF1set$)
 apply (erule UnE)
 apply (erule UN_E)
 apply (erule $F1set2_JF1set_JF1set$)
 apply assumption
 apply (erule UN_E)
 apply (erule $F1set3_JF2set_JF1set$)
 apply assumption
 apply (rule $subsetI$)
 apply (rule $UnI2$)
 apply (tactic $\langle\langle rtac\ @\{context\}\ (BNF_Util.mk_UnIN\ 2\ 2)\ 1\ \rangle\rangle$)
 apply (erule UN_I)
 apply (erule UnE)
 apply (erule $F2set1_JF2set$)
 apply (erule UnE)
 apply (erule UN_E)
 apply (erule $F2set2_JF1set_JF2set$)
 apply assumption
 apply (erule UN_E)
 apply (erule $F2set3_JF2set_JF2set$)
 apply assumption
 apply (rule $subsetI$)
 apply (rule $UnI2$)
 apply (tactic $\langle\langle rtac\ @\{context\}\ (BNF_Util.mk_UnIN\ 2\ 1)\ 1\ \rangle\rangle$)
 apply (erule UN_I)
 apply (erule UnE)
 apply (erule $F1set1_JF1set$)
 apply (erule UnE)
 apply (erule UN_E)
 apply (erule $F1set2_JF1set_JF1set$)
 apply assumption
 apply (erule UN_E)
 apply (erule $F1set3_JF2set_JF1set$)
 apply assumption
 apply (rule $subsetI$)
 apply (rule $UnI2$)
 apply (tactic $\langle\langle rtac\ @\{context\}\ (BNF_Util.mk_UnIN\ 2\ 2)\ 1\ \rangle\rangle$)
 apply (erule UN_I)
 apply (erule UnE)
 apply (erule $F2set1_JF2set$)
 apply (erule UnE)

```

apply (erule UN_E)
apply (erule F2set2_JF1set_JF2set)
apply assumption
apply (erule UN_E)
apply (erule F2set3_JF2set_JF2set)
apply assumption
done

```

theorem *JF1set_simps*:

```

JF1set a = F1set1 (dtor1 a) ∪
  ((∪ b ∈ F1set2 (dtor1 a). JF1set b) ∪
   (∪ b ∈ F1set3 (dtor1 a). JF2set b))
apply (rule equalityI)
apply (rule conjunct1[OF JFset_le])
apply (rule Un_least)
apply (rule F1set1_incl_JF1set)
apply (rule Un_least)
apply (rule UN_least)
apply (erule F1set2_JF1set_incl_JF1set)
apply (rule UN_least)
apply (erule F1set3_JF2set_incl_JF1set)
done

```

theorem *JF2set_simps*:

```

JF2set a = F2set1 (dtor2 a) ∪
  ((∪ b ∈ F2set2 (dtor2 a). JF1set b) ∪
   (∪ b ∈ F2set3 (dtor2 a). JF2set b))
apply (rule equalityI)
apply (rule conjunct2[OF JFset_le])
apply (rule Un_least)
apply (rule F2set1_incl_JF2set)
apply (rule Un_least)
apply (rule UN_least)
apply (erule F2set2_JF1set_incl_JF2set)
apply (rule UN_least)
apply (erule F2set3_JF2set_incl_JF2set)
done

```

lemma *JFcol_natural*:

```

∀ b1 b2. u ‘ (JF1col n b1) = JF1col n (JF1map u b1) ∧
  u ‘ (JF2col n b2) = JF2col n (JF2map u b2)
apply (rule nat_induct)
apply (rule allI)+
unfolding JF1col_0 JF2col_0
apply (rule conjI)
apply (rule image_empty)
apply (rule image_empty)

apply (rule allI)+
apply (rule conjI)
apply (unfold JF1col_Suc JF1map_simps image_Un image_UN UN_simps(10)
  F1.set_map(1) F1.set_map(2) F1.set_map(3)) [1]
apply (rule arg_cong2[of _ _ _ (∪)])
apply (rule refl)
apply (rule arg_cong2[of _ _ _ (∪)])
apply (rule SUP_cong[OF refl])
apply (erule allE)+
apply (tactic ⟨ etac @{context} (BNF_Util.mk_conjunctN 2 1) 1 ⟩)
apply (rule SUP_cong[OF refl])
apply (erule allE)+
apply (tactic ⟨ etac @{context} (BNF_Util.mk_conjunctN 2 2) 1 ⟩)

apply (unfold JF2col_Suc JF2map_simps image_Un image_UN UN_simps(10))

```

```

    F2.set_map(1) F2.set_map(2) F2.set_map(3)) [1]
apply (rule arg_cong2[of _ _ _ _ (U)])
apply (rule refl)
apply (rule arg_cong2[of _ _ _ _ (U)])
apply (rule SUP_cong[OF refl])
apply (erule allE)+
apply (tactic << etac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (rule SUP_cong[OF refl])
apply (erule allE)+
apply (tactic << etac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
done

```

theorem JF1set_natural: $JF1set \circ (JF1map \ u) = image \ u \circ JF1set$

```

apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule sym)
apply (rule trans[OF o_apply])
apply (rule trans[OF image_UN])
apply (rule SUP_cong[OF refl])
apply (rule conjunct1)
apply (rule spec[OF spec[OF JFcol_natural]])
done

```

theorem JF2set_natural: $JF2set \circ (JF2map \ u) = image \ u \circ JF2set$

```

apply (rule ext)
apply (rule trans[OF o_apply])
apply (rule sym)
apply (rule trans[OF o_apply])
apply (rule trans[OF image_UN])
apply (rule SUP_cong[OF refl])
apply (rule conjunct2)
apply (rule spec[OF spec[OF JFcol_natural]])
done

```

theorem JFmap_cong0:

```

(( $\forall p \in JF1set \ a. \ u \ p = v \ p$ )  $\longrightarrow$   $JF1map \ u \ a = JF1map \ v \ a$ )  $\wedge$ 
(( $\forall p \in JF2set \ b. \ u \ p = v \ p$ )  $\longrightarrow$   $JF2map \ u \ b = JF2map \ v \ b$ )
apply (rule rev_mp)
apply (rule Frel_coind[of
  %b c.  $\exists a. a \in \{a. \forall p \in JF1set \ a. \ u \ p = v \ p\} \wedge b = JF1map \ u \ a \wedge c = JF1map \ v \ a$ 
  %b c.  $\exists a. a \in \{a. \forall p \in JF2set \ a. \ u \ p = v \ p\} \wedge b = JF2map \ u \ a \wedge c = JF2map \ v \ a$ ])
apply (intro allI impI iffD2[OF F1.in_rel])

```

```

apply (erule exE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule exI)

```

```

apply (rule conjI[rotated])
apply (rule conjI)
apply (rule trans[OF F1.map_comp])
apply (rule sym)
apply (rule trans[OF JF1map_simps])
apply (rule F1.map_cong0)
apply (rule sym[OF trans[OF o_apply]])
apply (rule fst_conv)
apply (rule sym[OF fun_cong[OF fst_conv[unfolded_conv_def]])]
apply (rule sym[OF fun_cong[OF fst_conv[unfolded_conv_def]])]

```

```

apply (rule trans[OF F1.map_comp])
apply (rule sym)
apply (rule trans[OF JF1map_simps])
apply (rule F1.map_cong0)
apply (rule sym[OF trans[OF o_apply]])

```

```

apply (rule trans[OF snd_conv])
apply (erule CollectE)
apply (erule bspec)
apply (erule F1set1_JF1set)
apply (rule sym[OF fun_cong[OF snd_conv[unfolded convol_def]]])
apply (rule sym[OF fun_cong[OF snd_conv[unfolded convol_def]])

```

```

apply (rule CollectI)
apply (rule conjI)
apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(1))
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule case_prodI)
apply (rule refl)

```

```

apply (rule conjI)
apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(2))
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule case_prodI)
apply (rule exI)
apply (rule conjI)
apply (rule CollectI)
apply (rule ballI)
apply (erule CollectE)
apply (erule bspec)
apply (erule F1set2_JF1set_JF1set)
apply assumption
apply (rule conjI[OF refl refl])

```

```

apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(3))
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule case_prodI)
apply (rule exI)
apply (rule conjI)
apply (rule CollectI)
apply (rule ballI)
apply (erule CollectE)
apply (erule bspec)
apply (erule F1set3_JF2set_JF1set)
apply assumption
apply (rule conjI[OF refl refl])

```

```

apply (intro allI impI iffD2[OF F2.in_rel])

```

```

apply (erule exE conjE)+
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (rule exI)

```

```

apply (rule conjI[rotated])
apply (rule conjI)
apply (rule trans[OF F2.map_comp])
apply (rule sym)
apply (rule trans[OF JF2map_simps])
apply (rule F2.map_cong0)
  apply (rule sym[OF trans[OF o_apply]])

```

```

    apply (rule fst_conv)
    apply (rule sym[OF fun_cong[OF fst_conv[unfolded convol_def]])]
    apply (rule sym[OF fun_cong[OF fst_conv[unfolded convol_def]])]

apply (rule trans[OF F2.map_comp])
apply (rule sym)
apply (rule trans[OF JF2map_simps])
apply (rule F2.map_cong0)
  apply (rule sym[OF trans[OF o_apply]])
  apply (rule trans[OF snd_conv])
  apply (erule CollectE)
  apply (erule bspec)
  apply (erule F2set1_JF2set)
  apply (rule sym[OF fun_cong[OF snd_conv[unfolded convol_def]])]
  apply (rule sym[OF fun_cong[OF snd_conv[unfolded convol_def]])]

apply (rule CollectI)
apply (rule conjI)
  apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(1))
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule case_prodI)
  apply (rule refl)

apply (rule conjI)
  apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(2))
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule case_prodI)
  apply (rule exI)
  apply (rule conjI)
  apply (rule CollectI)
  apply (rule ballI)
  apply (erule CollectE)
  apply (erule bspec)
  apply (erule F2set2_JF1set_JF2set)
  apply assumption
  apply (rule conjI[OF refl refl])

apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(3))
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule case_prodI)
  apply (rule exI)
  apply (rule conjI)
  apply (rule CollectI)
  apply (rule ballI)
  apply (erule CollectE)
  apply (erule bspec)
  apply (erule F2set3_JF2set_JF2set)
  apply assumption
  apply (rule conjI[OF refl refl])

apply (rule impI)

apply (rule conjI)
  apply (rule impI)
  apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)

```

```

apply (erule mp)
apply (rule exI)
apply (rule conjI)
  apply (erule CollectI)
apply (rule conjI)
  apply (rule refl)
apply (rule refl)

apply (rule impI)
apply (tactic << dtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (erule mp)
apply (rule exI)
apply (rule conjI)
  apply (erule CollectI)
apply (rule conjI)
  apply (rule refl)
apply (rule refl)
done

lemmas JF1map_cong0 = mp[OF conjunct1[OF JFmap_cong0]]
lemmas JF2map_cong0 = mp[OF conjunct2[OF JFmap_cong0]]

lemma JFcol_bd:  $\forall (j1 :: 'a JF1) (j2 :: 'a JF2). |JF1col\ n\ j1| \leq o\ bd\_F \wedge |JF2col\ n\ j2| \leq o\ bd\_F$ 
apply (rule nat_induct)
apply (rule allI)+
apply (rule conjI)
  apply (rule ordIso_ordLeq_trans)
    apply (rule card_of_ordIso_subst)
      apply (rule JF1col_0)
    apply (rule Card_order_empty)
    apply (rule bd_F_Card_order)
  apply (rule ordIso_ordLeq_trans)
    apply (rule card_of_ordIso_subst)
      apply (rule JF2col_0)
    apply (rule Card_order_empty)
    apply (rule bd_F_Card_order)

apply (rule allI)+
apply (rule conjI)
  apply (rule ordIso_ordLeq_trans)
    apply (rule card_of_ordIso_subst)
      apply (rule JF1col_Suc)
  apply (rule Un_Cinfinite_bound)
    apply (rule F1set1_bd')
  apply (rule Un_Cinfinite_bound)
    apply (rule UNION_Cinfinite_bound)
      apply (rule F1set2_bd')
      apply (rule ballI)
      apply (erule allE)+
      apply (tactic << etac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
      apply (rule bd_F_Cinfinite)
  apply (rule UNION_Cinfinite_bound)
    apply (rule F1set3_bd')
    apply (rule ballI)
    apply (erule allE)+
    apply (tactic << etac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
    apply (rule bd_F_Cinfinite)
  apply (rule bd_F_Cinfinite)
apply (rule bd_F_Cinfinite)

apply (rule ordIso_ordLeq_trans)
apply (rule card_of_ordIso_subst)
apply (rule JF2col_Suc)

```

```

apply (rule Un_Cinfinite_bound)
apply (rule F2set1_bd')
apply (rule Un_Cinfinite_bound)
apply (rule UNION_Cinfinite_bound)
apply (rule F2set2_bd')
apply (rule ballI)
apply (erule allE)+
apply (tactic << etac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (rule bd_F_Cinfinite)
apply (rule UNION_Cinfinite_bound)
apply (rule F2set3_bd')
apply (rule ballI)
apply (erule allE)+
apply (tactic << etac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (rule bd_F_Cinfinite)
apply (rule bd_F_Cinfinite)
apply (rule bd_F_Cinfinite)
done

```

```

theorem JF1set_bd: |JF1set j| ≤o bd_F
apply (rule UNION_Cinfinite_bound)
apply (rule ordIso_ordLeq_trans)
apply (rule card_of_nat)
apply (rule natLeq_ordLeq_cinfinite)
apply (rule bd_F_Cinfinite)
apply (rule ballI)
apply (tactic << rtac @{context} (BNF_Util.mk_conjunctN 2 1) 1 >>)
apply (rule spec[OF spec[OF JFcol_bd]])
apply (rule bd_F_Cinfinite)
done

```

```

theorem JF2set_bd: |JF2set j| ≤o bd_F
apply (rule UNION_Cinfinite_bound)
apply (rule ordIso_ordLeq_trans)
apply (rule card_of_nat)
apply (rule natLeq_ordLeq_cinfinite)
apply (rule bd_F_Cinfinite)
apply (rule ballI)
apply (tactic << rtac @{context} (BNF_Util.mk_conjunctN 2 2) 1 >>)
apply (rule spec[OF spec[OF JFcol_bd]])
apply (rule bd_F_Cinfinite)
done

```

abbreviation JF2wit ≡ ctor2 wit_F2

```

theorem JF2wit: ∧x. x ∈ JF2set JF2wit ⇒ False
apply (drule set_rev_mp)
apply (rule equalityD1)
apply (rule JF2set_simps)
unfolding dtor2_ctor2
apply (erule UnE)
apply (erule F2.wit)
apply (erule UnE)
apply (erule UN_E)
apply (erule F2.wit)
apply (erule UN_E)
apply (erule F2.wit)
done

```

abbreviation JF1wit ≡ (%a. ctor1 (wit2_F1 a JF2wit))

```

theorem JF1wit: ∧x. x ∈ JF1set (JF1wit a) ⇒ x = a
apply (drule set_rev_mp)

```



```

  apply (rule equalityD1)
  apply (rule JF1set_simps)
  unfolding dtor1_ctor1
  apply (erule UnE)+
  apply (erule F1.wit2)
  apply (erule UnE)
  apply (erule UN_E)
  apply (erule F1.wit2)
  apply (erule FalseE)
  apply (erule UN_E)
  apply (erule F1.wit2)
  apply (tactic << hyp_subst_tac @{context} 1 >>)
  apply (erule set_rev_mp)
  apply (rule equalityD1)
  apply (rule JF2set_simps)
  unfolding dtor2_ctor2
  apply (erule UnE)+
  apply (erule F2.wit)
  apply (erule FalseE)
  apply (erule UnE)
  apply (erule UN_E)
  apply (erule F2.wit)
  apply (erule FalseE)
  apply (erule UN_E)
  apply (erule F2.wit)
  apply (erule FalseE)
done

```

context

fixes $\phi1 :: 'a \Rightarrow 'a \text{ JF1} \Rightarrow \text{bool}$ **and** $\phi2 :: 'a \Rightarrow 'a \text{ JF2} \Rightarrow \text{bool}$
begin

lemmas $JFset_induct =$

```

  JFset_minimal[of %b1. {a ∈ JF1set b1 . phi1 a b1} %b2. {a ∈ JF2set b2 . phi2 a b2},
  unfolded subset_Collect_iff[OF F1set1_incl_JF1set] subset_Collect_iff[OF F2set1_incl_JF2set]
  subset_Collect_iff[OF subset_refl],
  OF ballI ballI
  subset_CollectI[OF F1set2_JF1set_incl_JF1set]
  subset_CollectI[OF F1set3_JF2set_incl_JF1set]
  subset_CollectI[OF F2set2_JF1set_incl_JF2set]
  subset_CollectI[OF F2set3_JF2set_incl_JF2set]]

```

end

ML << rule_by_tactic @{context} (ALLGOALS (TRY o etac @{context} asm_rl)) @{thm JFset_induct} >>

abbreviation $JF1in$ **where** $JF1in B \equiv \{a. JF1set a \subseteq B\}$

abbreviation $JF2in$ **where** $JF2in B \equiv \{a. JF2set a \subseteq B\}$

definition $JF1rel$ **where**

```

  JF1rel R = (BNF_Def.Grp (JF1in (Collect (case_prod R))) (JF1map fst)) ^--1 OO
  (BNF_Def.Grp (JF1in (Collect (case_prod R))) (JF1map snd))

```

definition $JF2rel$ **where**

```

  JF2rel R = (BNF_Def.Grp (JF2in (Collect (case_prod R))) (JF2map fst)) ^--1 OO
  (BNF_Def.Grp (JF2in (Collect (case_prod R))) (JF2map snd))

```

lemma *in_JF1rel*:

$JF1rel\ R\ x\ y \longleftrightarrow (\exists z. z \in JF1in\ (Collect\ (case_prod\ R)) \wedge JF1map\ fst\ z = x \wedge JF1map\ snd\ z = y)$
by (*rule predicate2_eqD[OF trans[OF JF1rel_def OO_Grp_alt]]*)

lemma *in_JF2rel*:

$JF2rel\ R\ x\ y \longleftrightarrow (\exists z. z \in JF2in\ (Collect\ (case_prod\ R)) \wedge JF2map\ fst\ z = x \wedge JF2map\ snd\ z = y)$
by (*rule predicate2_eqD[OF trans[OF JF2rel_def OO_Grp_alt]]*)

lemma *J_rel_coind_ind*:

$\llbracket \forall x\ y. R2\ x\ y \longrightarrow (f\ x\ y \in F1in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))) \wedge$
 $F1map\ fst\ fst\ fst\ (f\ x\ y) = dtor1\ x \wedge$
 $F1map\ snd\ snd\ snd\ (f\ x\ y) = dtor1\ y;$
 $\forall x\ y. R3\ x\ y \longrightarrow (g\ x\ y \in F2in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))) \wedge$
 $F2map\ fst\ fst\ fst\ (g\ x\ y) = dtor2\ x \wedge$
 $F2map\ snd\ snd\ snd\ (g\ x\ y) = dtor2\ y \rrbracket \Longrightarrow$

$(\forall a \in JF1set\ z1. \forall x\ y. R2\ x\ y \wedge z1 = unfold1\ (case_prod\ f)\ (case_prod\ g)\ (x, y) \longrightarrow R1\ (fst\ a)\ (snd\ a)) \wedge$
 $(\forall a \in JF2set\ z2. \forall x\ y. R3\ x\ y \wedge z2 = unfold2\ (case_prod\ f)\ (case_prod\ g)\ (x, y) \longrightarrow R1\ (fst\ a)\ (snd\ a))$

apply (*tactic* $\llbracket rtac\ @\{context\}\ (rule_by_tactic\ @\{context\}\ (ALLGOALS\ (TRY\ o\ etac\ @\{context\}\ asm_rl))$
 $@\{thm\ JFset_induct\ of$

$\lambda a\ z1. \forall x\ y. R2\ x\ y \wedge z1 = unfold1\ (case_prod\ f)\ (case_prod\ g)\ (x, y) \longrightarrow R1\ (fst\ a)\ (snd\ a)$

$\lambda a\ z2. \forall x\ y. R3\ x\ y \wedge z2 = unfold2\ (case_prod\ f)\ (case_prod\ g)\ (x, y) \longrightarrow R1\ (fst\ a)\ (snd\ a)$

$z1\ z2\} 1 \rrbracket$)

apply (*rule allI impI*)+

apply (*erule conjE*)

apply (*drule spec2*)

apply (*erule thin_rl*)

apply (*drule mp*)

apply *assumption*

apply (*erule CollectE conjE Collect_case_prodD[OF set_mp] set_rev_mp*)+

apply *hypsubst*

unfolding *unfold1 F1.set_map(1) prod.case_image_id id_apply*

apply (*rule subset_refl*)

apply (*rule allI impI*)+

apply (*erule conjE*)

apply (*erule thin_rl*)

apply (*drule spec2*)

apply (*drule mp*)

apply *assumption*

apply (*erule CollectE conjE Collect_case_prodD[OF set_mp] set_rev_mp*)+

apply *hypsubst*

unfolding *unfold2 F2.set_map(1) prod.case_image_id id_apply*

apply (*rule subset_refl*)

apply (*rule impI allI*)+

apply (*erule conjE*)

apply (*drule spec2*)

apply (*erule thin_rl*)

apply (*drule mp*)

apply *assumption*

apply (*erule CollectE conjE*)+

apply (*tactic* $\llbracket hyp_subst_tac\ @\{context\}\ 1 \rrbracket$)

unfolding *unfold1 F1.set_map(2) prod.case_image_id id_apply*

apply (*erule imageE*)

apply (*tactic* $\llbracket hyp_subst_tac\ @\{context\}\ 1 \rrbracket$)

apply (*erule allE mp*)+

apply (*rule conjI*)

apply (*erule Collect_case_prodD[OF set_mp]*)

apply *assumption*

apply (*rule arg_cong[OF surjective_pairing]*)

```

apply (rule impI allI)+
apply (erule conjE)
apply (drule spec2)
apply (erule thin_rl)
apply (drule mp)
  apply assumption
apply (erule CollectE conjE)+
apply (tactic << hyp_subst_tac @ {context} 1 >>)
unfolding unfold1 F1.set_map(3) prod.case image_id id_id apply
apply (erule imageE)
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (erule allE mp)+
apply (rule conjI)
  apply (erule Collect_case_prodD[OF set_mp])
  apply assumption
apply (rule arg_cong[OF surjective_pairing])

```

```

apply (rule impI allI)+
apply (erule conjE)
apply (erule thin_rl)
apply (drule spec2)
apply (drule mp)
  apply assumption
apply (erule CollectE conjE)+
apply (tactic << hyp_subst_tac @ {context} 1 >>)
unfolding unfold2 F2.set_map(2) prod.case image_id id_id apply
apply (erule imageE)
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (erule allE mp)+
apply (rule conjI)
  apply (erule Collect_case_prodD[OF set_mp])
  apply assumption
apply (rule arg_cong[OF surjective_pairing])

```

```

apply (rule impI allI)+
apply (erule conjE)
apply (erule thin_rl)
apply (drule spec2)
apply (drule mp)
  apply assumption
apply (erule CollectE conjE)+
apply (tactic << hyp_subst_tac @ {context} 1 >>)
unfolding unfold2 F2.set_map(3) prod.case image_id id_id apply
apply (erule imageE)
apply (tactic << hyp_subst_tac @ {context} 1 >>)
apply (erule allE mp)+
apply (rule conjI)
  apply (erule Collect_case_prodD[OF set_mp])
  apply assumption
apply (rule arg_cong[OF surjective_pairing])
done

```

lemma J_rel_coind_coind1:

$$\begin{aligned}
& \llbracket \forall x y. R2 \ x \ y \longrightarrow (f \ x \ y \in F1in \ (Collect \ (case_prod \ R1)) \ (Collect \ (case_prod \ R2)) \ (Collect \ (case_prod \ R3))) \wedge \\
& \quad F1map \ fst \ fst \ fst \ (f \ x \ y) = dtor1 \ x \wedge \\
& \quad F1map \ snd \ snd \ snd \ (f \ x \ y) = dtor1 \ y; \\
& \forall x y. R3 \ x \ y \longrightarrow (g \ x \ y \in F2in \ (Collect \ (case_prod \ R1)) \ (Collect \ (case_prod \ R2)) \ (Collect \ (case_prod \ R3))) \wedge \\
& \quad F2map \ fst \ fst \ fst \ (g \ x \ y) = dtor2 \ x \wedge
\end{aligned}$$

```

      F2map snd snd snd (g x y) = dtor2 y]] ==>
((∃ y. R2 x1 y ∧ x1' = JF1map fst (unfold1 (case_prod f) (case_prod g) (x1, y))) → x1' = x1) ∧
((∃ y. R3 x2 y ∧ x2' = JF2map fst (unfold2 (case_prod f) (case_prod g) (x2, y))) → x2' = x2)
apply (rule Frel_coind[of
  λx1' x1. ∃ y. R2 x1 y ∧ x1' = JF1map fst (unfold1 (case_prod f) (case_prod g) (x1, y))
  λx2' x2. ∃ y. R3 x2 y ∧ x2' = JF2map fst (unfold2 (case_prod f) (case_prod g) (x2, y))
  x1' x1
  x2' x2
])
apply (intro allI impI iffD2[OF F1.in_rel])

apply (erule exE conjE)+
apply (drule spec2)
apply (erule thin_rl)
apply (drule mp)
apply assumption
apply (erule CollectE conjE)+
apply (tactic ⟨ hyp_subst_tac @ {context} 1 ⟩)
apply (rule exI)
apply (rule conjI[rotated])
apply (rule conjI)
apply (rule trans[OF F1.map_comp])
apply (rule sym[OF trans[OF JF1map_simps]])
apply (rule trans[OF arg_cong[OF unfold1]])
apply (rule trans[OF F1.map_comp F1.map_cong0])
apply (rule trans[OF fun_cong[OF o_id]])
apply (rule sym[OF fun_cong[OF fst_diag_fst]])
apply (rule sym[OF trans[OF o_apply]])
apply (rule fst_conv)
apply (rule sym[OF trans[OF o_apply]])
apply (rule fst_conv)
apply (rule trans[OF F1.map_comp])
apply (rule trans[OF F1.map_cong0])
apply (rule fun_cong[OF snd_diag_fst])
apply (rule trans[OF o_apply])
apply (rule snd_conv)
apply (rule trans[OF o_apply])
apply (rule snd_conv)
apply (erule trans[OF arg_cong[OF prod.case]])

apply (unfold prod.case o_def fst_conv snd_conv) []
apply (rule CollectI)
apply (rule conjI)
apply (rule ord_eq_le_trans[OF F1.set_map(1)])
apply (rule image_subsetI CollectI case_prodI)+
apply (rule refl)

apply (rule conjI)
apply (rule ord_eq_le_trans[OF F1.set_map(2)])
apply (rule image_subsetI CollectI case_prodI exI)+
apply (rule conjI)
apply (erule Collect_case_prodD[OF set_mp])
apply assumption
apply (rule arg_cong[OF surjective_pairing])

apply (rule ord_eq_le_trans[OF F1.set_map(3)])
apply (rule image_subsetI CollectI case_prodI exI)+
apply (rule conjI)
apply (erule Collect_case_prodD[OF set_mp])
apply assumption
apply (rule arg_cong[OF surjective_pairing])

```

```

apply (intro allI impI iffD2[OF F2.in_rel])

apply (erule exE conjE)+
apply (erule thin_rl)
apply (drule spec2)
apply (drule mp)
apply assumption
apply (erule CollectE conjE)+
apply (tactic ⟨⟨ hyp_subst_tac @ {context} 1 ⟩⟩)
apply (rule exI)
apply (rule conjI[rotated])
apply (rule conjI)
apply (rule trans[OF F2.map_comp])
apply (rule sym[OF trans[OF JF2map_simps]])
apply (rule trans[OF arg_cong[OF unfold2]])
apply (rule trans[OF F2.map_comp F2.map_cong0])
apply (rule fun_cong[OF trans[OF o_id fst_diag_fst[symmetric]]])
apply (rule sym[OF trans[OF o_apply]])
apply (rule fst_conv)
apply (rule sym[OF trans[OF o_apply]])
apply (rule fst_conv)
apply (rule trans[OF F2.map_comp])
apply (rule trans[OF F2.map_cong0])
apply (rule fun_cong[OF snd_diag_fst])
apply (rule trans[OF o_apply])
apply (rule snd_conv)
apply (rule trans[OF o_apply])
apply (rule snd_conv)
apply (erule trans[OF arg_cong[OF prod.case]])

apply (unfold prod.case o_def fst_conv snd_conv) []
apply (rule CollectI)
apply (rule conjI)
apply (rule ord_eq_le_trans[OF F2.set_map(1)])
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule case_prodI)
apply (rule refl)

apply (rule conjI)
apply (rule ord_eq_le_trans[OF F2.set_map(2)])
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule case_prodI)
apply (rule exI)
apply (rule conjI)
apply (erule Collect_case_prodD[OF set_mp])
apply assumption
apply (rule arg_cong[OF surjective_pairing])

apply (rule ord_eq_le_trans[OF F2.set_map(3)])
apply (rule image_subsetI CollectI case_prodI exI)+
apply (rule conjI)
apply (erule Collect_case_prodD[OF set_mp])
apply assumption
apply (rule arg_cong[OF surjective_pairing])
done

```

lemma *J_rel_coind_coind2*:

$$\llbracket \forall x y. R2 \ x \ y \longrightarrow (f \ x \ y \in F1in \ (Collect \ (case_prod \ R1)) \ (Collect \ (case_prod \ R2)) \ (Collect \ (case_prod \ R3))) \wedge \\ F1map \ fst \ fst \ fst \ (f \ x \ y) = dtor1 \ x \wedge \\ F1map \ snd \ snd \ snd \ (f \ x \ y) = dtor1 \ y \rrbracket;$$

```

 $\forall x y. R3 x y \longrightarrow (g x y \in F2in (Collect (case\_prod R1)) (Collect (case\_prod R2)) (Collect (case\_prod R3)) \wedge$ 
 $F2map fst fst fst (g x y) = dtor2 x \wedge$ 
 $F2map snd snd snd (g x y) = dtor2 y) \implies$ 
 $((\exists x. R2 x y1 \wedge y1' = JF1map snd (unfold1 (case\_prod f) (case\_prod g) (x, y1))) \longrightarrow y1' = y1) \wedge$ 
 $((\exists x. R3 x y2 \wedge y2' = JF2map snd (unfold2 (case\_prod f) (case\_prod g) (x, y2))) \longrightarrow y2' = y2)$ 
apply (rule Frel_coind[of
   $\lambda y1' y1. \exists x. R2 x y1 \wedge y1' = JF1map snd (unfold1 (case\_prod f) (case\_prod g) (x, y1))$ 
   $\lambda y2' y2. \exists x. R3 x y2 \wedge y2' = JF2map snd (unfold2 (case\_prod f) (case\_prod g) (x, y2))$ 
   $y1' y1$ 
   $y2' y2$ 
])
apply (intro allI impI iffD2[OF F1.in_rel])

apply (erule exE conjE)+
apply (erule spec2)
apply (erule thin_rl)
apply (erule mp)
apply assumption
apply (erule CollectE conjE)+
apply (tactic { hyp_subst_tac @ {context} 1 })
apply (rule exI)
apply (rule conjI[rotated])
apply (rule conjI)
apply (rule trans[OF F1.map_comp])
apply (rule sym[OF trans[OF JF1map_simps]])
apply (rule trans[OF arg_cong[OF unfold1]])
apply (rule trans[OF F1.map_comp F1.map_cong0])
apply (rule trans[OF fun_cong[OF o_id]])
apply (rule sym[OF fun_cong[OF fst_diag_snd]])
apply (rule sym[OF trans[OF o_apply]])
apply (rule fst_conv)
apply (rule sym[OF trans[OF o_apply]])
apply (rule fst_conv)
apply (rule trans[OF F1.map_comp])
apply (rule trans[OF F1.map_cong0])
apply (rule fun_cong[OF snd_diag_snd])
apply (rule trans[OF o_apply])
apply (rule snd_conv)
apply (rule trans[OF o_apply])
apply (rule snd_conv)
apply (erule trans[OF arg_cong[OF prod.case]])

apply (unfold prod.case o_def fst_conv snd_conv) []
apply (rule CollectI)
apply (rule conjI)
apply (rule ord_eq_le_trans[OF F1.set_map(1)])
apply (rule image_subsetI CollectI case_prodI)+
apply (rule refl)

apply (rule conjI)
apply (rule ord_eq_le_trans[OF F1.set_map(2)])
apply (rule image_subsetI CollectI case_prodI exI)+
apply (rule conjI)
apply (erule Collect_case_prodD[OF set_mp])
apply assumption
apply (rule arg_cong[OF surjective_pairing])

apply (rule ord_eq_le_trans[OF F1.set_map(3)])
apply (rule image_subsetI CollectI case_prodI exI)+
apply (rule conjI)
apply (erule Collect_case_prodD[OF set_mp])
apply assumption
apply (rule arg_cong[OF surjective_pairing])

```

```

apply (intro allI impI iffD2[OF F2.in_rel])

apply (erule exE conjE)+
apply (erule thin_rl)
apply (drule spec2)
apply (drule mp)
  apply assumption
apply (erule CollectE conjE)+
apply (tactic << hyp_subst_tac @{context} 1 >>)
apply (rule exI)
apply (rule conjI[rotated])
apply (rule conjI)
  apply (rule trans[OF F2.map_comp])
  apply (rule sym[OF trans[OF JF2map_simps]])
  apply (rule trans[OF arg_cong[OF unfold2]])
apply (rule trans[OF F2.map_comp F2.map_cong0])
  apply (rule trans[OF fun_cong[OF o_id]])
  apply (rule sym[OF fun_cong[OF fst_diag_snd]])
  apply (rule sym[OF trans[OF o_apply]])
  apply (rule fst_conv)
apply (rule sym[OF trans[OF o_apply]])
apply (rule fst_conv)
apply (rule trans[OF F2.map_comp])
apply (rule trans[OF F2.map_cong0])
  apply (rule fun_cong[OF snd_diag_snd])
  apply (rule trans[OF o_apply])
  apply (rule snd_conv)
apply (rule trans[OF o_apply])
apply (rule snd_conv)
apply (erule trans[OF arg_cong[OF prod.case]])

apply (unfold prod.case o_def fst_conv snd_conv) []
apply (rule CollectI)
apply (rule conjI)
  apply (rule ord_eq_le_trans[OF F2.set_map(1)])
  apply (rule image_subsetI CollectI case_prodI)+
apply (rule refl)

apply (rule conjI)
  apply (rule ord_eq_le_trans[OF F2.set_map(2)])
  apply (rule image_subsetI CollectI case_prodI exI)+
apply (rule conjI)
  apply (erule Collect_case_prodD[OF set_mp])
  apply assumption
apply (rule arg_cong[OF surjective_pairing])

apply (rule ord_eq_le_trans[OF F2.set_map(3)])
apply (rule image_subsetI CollectI case_prodI exI)+
apply (rule conjI)
  apply (erule Collect_case_prodD[OF set_mp])
  apply assumption
apply (rule arg_cong[OF surjective_pairing])
done

lemma J_rel_coind:
assumes CIH1:  $\forall x2 y2. R2 x2 y2 \longrightarrow F1rel R1 R2 R3 (dtor1 x2) (dtor1 y2)$ 
  and CIH2:  $\forall x3 y3. R3 x3 y3 \longrightarrow F2rel R1 R2 R3 (dtor2 x3) (dtor2 y3)$ 
shows  $R2 \leq JF1rel R1 \wedge R3 \leq JF2rel R1$ 
apply (insert CIH1 CIH2)
unfolding F1.in_rel F2.in_rel ex_simps(6)[symmetric] choice_iff

```

```

apply (erule exE)+
apply (rule conjI)

apply (rule predicate2I)
apply (rule iffD2[OF in_JF1rel])
apply (rule exI conjI)+
apply (rule CollectI)
apply (rule rev_mp[OF conjunct1[OF J_rel_coind_ind]])
  apply assumption
  apply assumption
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule impI)
apply (rule subsetI CollectI iffD2[OF case_prod_beta])+
apply (drule bspec)
  apply assumption
apply (erule allE mp conjE)+
apply (erule conjI[OF _ refl])

apply (rule conjI)
apply (rule rev_mp[OF conjunct1[OF J_rel_coind_coind1]])
  apply assumption
  apply assumption
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule impI)
apply (erule mp)
apply (rule exI)
apply (erule conjI[OF _ refl])

apply (rule rev_mp[OF conjunct1[OF J_rel_coind_coind2]])
  apply assumption
  apply assumption
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule impI)
apply (erule mp)
apply (rule exI)
apply (erule conjI[OF _ refl])

apply (rule predicate2I)
apply (rule iffD2[OF in_JF2rel])
apply (rule exI conjI)+
apply (rule rev_mp[OF conjunct2[OF J_rel_coind_ind]])
  apply assumption
  apply assumption
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule impI)
apply (rule CollectI)
apply (rule subsetI CollectI iffD2[OF case_prod_beta])+
apply (drule bspec)
  apply assumption
apply (erule allE mp conjE)+
apply (erule conjI[OF _ refl])

apply (rule conjI)
apply (rule rev_mp[OF conjunct2[OF J_rel_coind_coind1]])
  apply assumption
  apply assumption
apply (erule thin_rl)

```



```

apply (erule thin_rl)
apply (rule impI)
apply (erule mp)
apply (rule exI)
apply (erule conjI[OF _ refl])

apply (rule rev_mp[OF conjunct2[OF J_rel_coind_coind2]])
  apply assumption
  apply assumption
apply (erule thin_rl)
apply (erule thin_rl)
apply (rule impI)
apply (erule mp)
apply (rule exI)
apply (erule conjI[OF _ refl])
done

lemma JF1rel_F1rel: JF1rel R a b  $\longleftrightarrow$  F1rel R (JF1rel R) (JF2rel R) (dtor1 a) (dtor1 b)
apply (rule iffI)
apply (erule iffD1[OF in_JF1rel])
apply (erule exE conjE CollectE)+
apply (rule iffD2[OF F1.in_rel])
apply (rule exI)
apply (rule conjI)
apply (rule CollectI)
apply (rule conjI)
  apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(1))
apply (rule ord_eq_le_trans)
  apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply (erule subset_trans[OF F1set1_incl_JF1set])

apply (rule conjI)
apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(2))
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule case_prodI)
apply (rule iffD2[OF in_JF1rel])
apply (rule exI)
apply (rule conjI)
  apply (rule CollectI)
  apply (erule subset_trans[OF F1set2_JF1set_incl_JF1set])
apply assumption
apply (rule conjI)
apply (rule refl)
apply (rule refl)

apply (rule ord_eq_le_trans)
  apply (rule F1.set_map(3))
apply (rule image_subsetI)
apply (rule CollectI)
apply (rule case_prodI)
apply (rule iffD2[OF in_JF2rel])
apply (rule exI)
apply (rule conjI)
  apply (rule CollectI)
  apply (erule subset_trans[OF F1set3_JF2set_incl_JF1set])
apply assumption
apply (rule conjI)
  apply (rule refl)
apply (rule refl)
apply (rule conjI)

```

```

apply (rule trans)
apply (rule F1.map_comp)
apply (rule trans)
apply (rule F1.map_cong0)
  apply (rule fun_cong[OF o_id])
apply (rule trans)
  apply (rule o_apply)
apply (rule fst_conv)
apply (rule trans)
  apply (rule o_apply)
apply (rule fst_conv)
apply (rule trans)
apply (rule sym)
apply (rule JF1map_simps)
apply (rule iffD2[OF dtor1_diff])
apply assumption

```

```

apply (rule trans)
apply (rule F1.map_comp)
apply (rule trans)
apply (rule F1.map_cong0)
  apply (rule fun_cong[OF o_id])
apply (rule trans)
  apply (rule o_apply)
apply (rule snd_conv)
apply (rule trans)
  apply (rule o_apply)
apply (rule snd_conv)
apply (rule trans)
apply (rule sym)
apply (rule JF1map_simps)
apply (rule iffD2[OF dtor1_diff])
apply assumption

```

```

apply (drule iffD1[OF F1.in_rel])
apply (erule exE conjE CollectE)+
apply (rule iffD2[OF in_JF1rel])
apply (rule exI)
apply (rule conjI)
apply (rule CollectI)
apply (rule ord_eq_le_trans)
apply (rule JF1set_simps)
apply (rule Un_least)
apply (rule ord_eq_le_trans)
apply (rule box_equals)
  apply (rule F1.set_map(1))
  apply (rule arg_cong[OF sym[OF dtor1_ctor1]])
  apply (rule trans[OF fun_cong[OF image_id] id_apply])
apply assumption
apply (rule Un_least)
apply (rule ord_eq_le_trans)
apply (rule SUP_cong[OF _ refl])
apply (rule box_equals[OF _ _ refl])
  apply (rule F1.set_map(2))
  apply (rule arg_cong[OF sym[OF dtor1_ctor1]])
apply (rule UN_least)
apply (drule set_rev_mp)
apply (erule image_mono)
apply (erule imageE)
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst

```

```

apply (drule iffD1[OF in_JF1rel])
apply (drule someI_ex)
apply (erule conjE)+
apply (erule CollectE conjE)+
apply assumption

apply (rule ord_eq_le_trans)
apply (rule trans[OF arg_cong[OF dtor1_ctor1]])
apply (rule arg_cong[OF F1.set_map(3)])
apply (rule UN_least)
apply (drule set_rev_mp)
apply (erule image_mono)
apply (erule imageE)
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (drule iffD1[OF in_JF2rel])
apply (drule someI_ex)
apply (erule exE conjE)+
apply (erule CollectD)

apply (rule conjI)

apply (rule iffD1[OF dtor1_diff])
apply (rule trans)
apply (rule JF1map_simps)
apply (rule box_equals)
apply (rule F1.map_comp)
apply (rule arg_cong[OF sym[OF dtor1_ctor1]])
apply (rule trans)
apply (rule F1.map_cong0)
apply (rule fun_cong[OF o_id])
apply (rule trans[OF o_apply])
apply (drule set_rev_mp)
apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (drule iffD1[OF in_JF1rel])
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply (rule trans[OF o_apply])
apply (drule set_rev_mp)
apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (drule iffD1[OF in_JF2rel])
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply assumption

apply (rule iffD1[OF dtor1_diff])
apply (rule trans)
apply (rule JF1map_simps)
apply (rule trans)
apply (rule arg_cong[OF dtor1_ctor1])
apply (rule trans)
apply (rule F1.map_comp)
apply (rule trans)
apply (rule F1.map_cong0)

```

```

  apply (rule fun_cong[OF o_id])
  apply (rule trans[OF o_apply])
  apply (drule set_rev_mp)
  apply assumption
  apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
  apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
  apply hypsubst
  apply (drule iffD1[OF in_JF1rel])
  apply (drule someI_ex)
  apply (erule conjE)+
  apply assumption
  apply (rule trans[OF o_apply])
  apply (drule set_rev_mp)
  apply assumption
  apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
  apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
  apply hypsubst
  apply (drule iffD1[OF in_JF2rel])
  apply (drule someI_ex)
  apply (erule conjE)+
  apply assumption
  apply assumption
done

```

lemma *JF2rel_F2rel*: $JF2rel\ R\ a\ b \iff F2rel\ R\ (JF1rel\ R)\ (JF2rel\ R)\ (dior2\ a)\ (dior2\ b)$

```

  apply (rule iffI)
  apply (drule iffD1[OF in_JF2rel])
  apply (erule exE conjE CollectE)+
  apply (rule iffD2[OF F2.in_rel])
  apply (rule exI)
  apply (rule conjI)
  apply (rule CollectI)
  apply (rule conjI)
  apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(1))
  apply (rule ord_eq_le_trans)
  apply (rule trans[OF fun_cong[OF image_id] id_apply])
  apply (rule subset_trans)
  apply (rule F2set1_incl_JF2set)
  apply assumption

  apply (rule conjI)
  apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(2))
  apply (rule image_subsetI)
  apply (rule CollectI)
  apply (rule case_prodI)
  apply (rule iffD2[OF in_JF1rel])
  apply (rule exI)
  apply (rule conjI)
  apply (rule CollectI)
  apply (rule subset_trans)
  apply (rule F2set2_JF1set_incl_JF2set)
  apply assumption
  apply assumption
  apply (rule conjI)
  apply (rule refl)
  apply (rule refl)

  apply (rule ord_eq_le_trans)
  apply (rule F2.set_map(3))
  apply (rule image_subsetI)
  apply (rule CollectI)

```

```

apply (rule case_prodI)
apply (rule iffD2[OF in_JF2rel])
apply (rule exI)
apply (rule conjI)
  apply (rule CollectI)
  apply (rule subset_trans)
    apply (rule F2set3_JF2set_incl_JF2set)
    apply assumption
  apply assumption
apply (rule conjI)
  apply (rule refl)
apply (rule refl)
apply (rule conjI)

```

```

apply (rule trans)
  apply (rule F2.map_comp)
apply (rule trans)
apply (rule F2.map_cong0)
  apply (rule fun_cong[OF o_id])
  apply (rule trans)
    apply (rule o_apply)
    apply (rule fst_conv)
apply (rule trans)
  apply (rule o_apply)
apply (rule fst_conv)
apply (rule trans)
  apply (rule sym)
  apply (rule JF2map_simps)
apply (rule iffD2)
  apply (rule dtor2_diff)
apply assumption

```

```

apply (rule trans)
apply (rule F2.map_comp)
apply (rule trans)
apply (rule F2.map_cong0)
  apply (rule fun_cong[OF o_id])
  apply (rule trans)
    apply (rule o_apply)
    apply (rule snd_conv)
apply (rule trans)
  apply (rule o_apply)
apply (rule snd_conv)
apply (rule trans)
apply (rule sym)
apply (rule JF2map_simps)
apply (rule iffD2)
apply (rule dtor2_diff)
apply assumption

```

```

apply (drule iffD1[OF F2.in_rel])
apply (erule exE conjE CollectE)+
apply (rule iffD2[OF in_JF2rel])
apply (rule exI)
apply (rule conjI)
apply (rule CollectI)
apply (rule ord_eq_le_trans)
apply (rule JF2set_simps)
apply (rule Un_least)
apply (rule ord_eq_le_trans)
  apply (rule trans)
  apply (rule trans)
  apply (rule arg_cong[OF dtor2_ctor2])

```

```

    apply (rule F2.set_map(1))
    apply (rule trans[OF fun_cong[OF image_id] id_apply])
  apply assumption
  apply (rule Un_least)
  apply (rule ord_eq_le_trans)
    apply (rule trans[OF arg_cong[OF dtor2_ctor2]])
    apply (rule arg_cong[OF F2.set_map(2)])
  apply (rule UN_least)
  apply (drule set_rev_mp)
  apply (erule image_mono)
  apply (erule imageE)
  apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
  apply (erule CollectE case_proxE iffD1[OF prod.inject, elim_format] conjE)+
  apply hypsubst
  apply (drule iffD1[OF in_JF1rel])
  apply (drule someI_ex)
  apply (erule conjE)+
  apply (erule CollectD)

  apply (rule ord_eq_le_trans)
  apply (rule trans[OF arg_cong[OF dtor2_ctor2]])
  apply (rule arg_cong[OF F2.set_map(3)])
  apply (rule UN_least)
  apply (drule set_rev_mp)
  apply (erule image_mono)
  apply (erule imageE)
  apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
  apply (erule CollectE case_proxE iffD1[OF prod.inject, elim_format] conjE)+
  apply hypsubst
  apply (drule iffD1[OF in_JF2rel])
  apply (drule someI_ex)
  apply (erule exE conjE)+
  apply (erule CollectD)

  apply (rule conjI)

  apply (rule iffD1)
  apply (rule dtor2_diff)
  apply (rule trans)
  apply (rule JF2map_simps)
  apply (rule trans)
  apply (rule arg_cong[OF dtor2_ctor2])
  apply (rule trans)
  apply (rule F2.map_comp)
  apply (rule trans)
  apply (rule F2.map_cong0)
    apply (rule fun_cong[OF o_id])
    apply (rule trans[OF o_apply])
    apply (drule set_rev_mp)
    apply assumption
    apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
    apply (erule CollectE case_proxE iffD1[OF prod.inject, elim_format] conjE)+
    apply hypsubst
    apply (drule iffD1[OF in_JF1rel])
    apply (drule someI_ex)
    apply (erule conjE)+
    apply assumption
  apply (rule trans[OF o_apply])
  apply (drule set_rev_mp)
  apply assumption
  apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
  apply (erule CollectE case_proxE iffD1[OF prod.inject, elim_format] conjE)+
  apply hypsubst

```

```

apply (drule iffD1[OF in_ JF2rel])
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply assumption

apply (rule iffD1)
apply (rule dtor2_diff)
apply (rule trans)
apply (rule JF2map_simps)
apply (rule trans)
apply (rule arg_cong[OF dtor2_ctor2])
apply (rule trans)
apply (rule F2.map_comp)
apply (rule trans)
apply (rule F2.map_cong0)
  apply (rule fun_cong[OF o_id])
apply (rule trans[OF o_apply])
apply (drule set_rev_mp)
  apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (drule iffD1[OF in_ JF1rel])
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply (rule trans[OF o_apply])
apply (drule set_rev_mp)
  apply assumption
apply (drule ssubst_mem[OF surjective_pairing[symmetric]])
apply (erule CollectE case_prodE iffD1[OF prod.inject, elim_format] conjE)+
apply hypsubst
apply (drule iffD1[OF in_ JF2rel])
apply (drule someI_ex)
apply (erule conjE)+
apply assumption
apply assumption
done

```

lemma JFrel_Comp_le:

$$JF1rel\ R1\ OO\ JF1rel\ R2 \leq JF1rel\ (R1\ OO\ R2) \wedge JF2rel\ R1\ OO\ JF2rel\ R2 \leq JF2rel\ (R1\ OO\ R2)$$

```

apply (rule J_rel_coind[OF allI[OF allI[OF impI]] allI[OF allI[OF impI]]])
apply (rule predicate2D[OF eq_refl[OF sym[OF F1.rel_compp]]])
apply (erule relcomppE)
apply (rule relcomppI)
apply (erule iffD1[OF JF1rel_F1rel])
apply (erule iffD1[OF JF1rel_F1rel])
apply (rule predicate2D[OF eq_refl[OF sym[OF F2.rel_compp]]])
apply (erule relcomppE)
apply (rule relcomppI)
apply (erule iffD1[OF JF2rel_F2rel])
apply (erule iffD1[OF JF2rel_F2rel])
done

```

context includes *lifting_syntax*

begin

lemma unfold_transfer:

$$((S \implies F1rel\ R\ S\ T) \implies (T \implies F2rel\ R\ S\ T) \implies S \implies JF1rel\ R) \text{ unfold1 unfold1 } \wedge$$

$$((S \implies F1rel\ R\ S\ T) \implies (T \implies F2rel\ R\ S\ T) \implies T \implies JF2rel\ R) \text{ unfold2 unfold2}$$

```

unfolding rel_fun_def_butlast_all_conj_distrib[symmetric] imp_conjR[symmetric]
unfolding rel_fun_iff_geq_image2p

```

```

apply (rule allI impI) +
apply (rule J_rel_coind)
apply (rule allI impI) +
apply (erule image2pE)
apply hypsubst
apply (unfold unfold1 unfold2) [1]
apply (rule rel_funD[OF rel_funD[OF rel_funD[OF rel_funD[OF F1.map_transfer]]]])
  apply (rule id_transfer)
  apply (rule rel_fun_image2p)
  apply (rule rel_fun_image2p)
apply (erule predicate2D)
apply (erule image2pI)

apply (rule allI impI) +
apply (erule image2pE)
apply hypsubst
apply (unfold unfold1 unfold2) [1]
apply (rule rel_funD[OF rel_funD[OF rel_funD[OF rel_funD[OF F2.map_transfer]]]])
  apply (rule id_transfer)
  apply (rule rel_fun_image2p)
  apply (rule rel_fun_image2p)
apply (erule predicate2D)
apply (erule image2pI)
done

end

ML ⟨⟨
  BNF_FP_Util.mk_xtor_co_iter_o_map_thms BNF_Util.Greatest_FP false 1 @{thm unfold_unique}
  @{thms JF1map JF2map} (map (BNF_Tactics.mk_pointfree2 @{context}) @{thms unfold1 unfold2})
  @{thms F1.map_comp0[symmetric] F2.map_comp0[symmetric]} @{thms F1.map_cong0 F2.map_cong0}
  ⟩⟩

ML ⟨⟨
  BNF_FP_Util.mk_xtor_co_iter_o_map_thms BNF_Util.Greatest_FP true 1 @{thm corec_unique}
  @{thms JF1map JF2map} (map (BNF_Tactics.mk_pointfree2 @{context}) @{thms corec1 corec2})
  @{thms F1.map_comp0[symmetric] F2.map_comp0[symmetric]} @{thms F1.map_cong0 F2.map_cong0}
  ⟩⟩

bnf 'a JF1
  map: JF1map
  sets: JF1set
  bd: bd_F
  wits: JF1wit
  rel: JF1rel
  apply –
  apply (rule JF1map_id)
  apply (rule JF1map_comp)
  apply (erule JF1map_cong0[OF ballI])
  apply (rule JF1set_natural)
  apply (rule bd_F_card_order)
  apply (rule conjunct1[OF bd_F_Cinfinite])
  apply (rule JF1set_bd)
  apply (rule conjunct1[OF JFrel_Comp_le])
  apply (rule JF1rel_def[unfolded OO_Grp_alt mem_Collect_eq])
  apply (erule JF1wit)
  done

bnf 'a JF2
  map: JF2map
  sets: JF2set
  bd: bd_F
  wits: JF2wit

```



```

rel: JF2rel
  apply -
    apply (rule JF2map_id)
    apply (rule JF2map_comp)
    apply (erule JF2map_cong0[OF ballI])
    apply (rule JF2set_natural)
    apply (rule bd_F_card_order)
    apply (rule conjunct1[OF bd_F_Cinfinite])
    apply (rule JF2set_bd)
    apply (rule conjunct2[OF JFrel_Comp_le])
    apply (rule JF2rel_def[unfolded OO_Grp_alt mem_Collect_eq])
    apply (erule JF2wit)
done

```

3 Normalized Composition of BNFs

Expected normal form: outer m-ary BNF is composed with m inner n-ary BNFs.

```

declare [[bnf_internals]]
bnf-axiomatization (dead 'p1, F1set1: 'a1, F1set2: 'a2) F1
  [wits: ('p1, 'a1, 'a2) F1]
  for map: F1map rel: F1rel
bnf-axiomatization (dead 'p2, F2set1: 'a1, F2set2: 'a2) F2
  [wits: 'a1  $\Rightarrow$  ('p2, 'a1, 'a2) F2 'a2  $\Rightarrow$  ('p2, 'a1, 'a2) F2]
  for map: F2map rel: F2rel
bnf-axiomatization (dead 'p3, F3set1: 'a1, F3set2: 'a2) F3
  [wits: 'a1  $\Rightarrow$  'a2  $\Rightarrow$  ('p3, 'a1, 'a2) F3]
  for map: F3map rel: F3rel
bnf-axiomatization (dead 'p, Gset1: 'b1, Gset2: 'b2, Gset3: 'b3) G
  [wits: 'b1  $\Rightarrow$  'b3  $\Rightarrow$  ('p, 'b1, 'b2, 'b3) G 'b2  $\Rightarrow$  'b3  $\Rightarrow$  ('p, 'b1, 'b2, 'b3) G]
  for map: Gmap rel: Grel
type-synonym ('p1, 'p2, 'p3, 'p, 'a1, 'a2) H =
  ('p, ('p1, 'a1, 'a2) F1, ('p2, 'a1, 'a2) F2, ('p3, 'a1, 'a2) F3) G
type-synonym ('p1, 'p2, 'p3, 'p) Hbd_type =
  ('p1 bd_type_F1 + 'p2 bd_type_F2 + 'p3 bd_type_F3)  $\times$  'p bd_type_G

abbreviation F1in where F1in A1 A2  $\equiv$  {x. F1set1 x  $\subseteq$  A1  $\wedge$  F1set2 x  $\subseteq$  A2}
abbreviation F2in where F2in A1 A2  $\equiv$  {x. F2set1 x  $\subseteq$  A1  $\wedge$  F2set2 x  $\subseteq$  A2}
abbreviation F3in where F3in A1 A2  $\equiv$  {x. F3set1 x  $\subseteq$  A1  $\wedge$  F3set2 x  $\subseteq$  A2}
abbreviation Gin where Gin A1 A2 A3  $\equiv$  {x. Gset1 x  $\subseteq$  A1  $\wedge$  Gset2 x  $\subseteq$  A2  $\wedge$  Gset3 x  $\subseteq$  A3}

abbreviation Gset where
  Gset  $\equiv$  BNF_Def.collect {Gset1, Gset2, Gset3}

abbreviation Hmap :: ('a1  $\Rightarrow$  'b1)  $\Rightarrow$  ('a2  $\Rightarrow$  'b2)  $\Rightarrow$ 
  ('p1, 'p2, 'p3, 'p, 'a1, 'a2) H  $\Rightarrow$  ('p1, 'p2, 'p3, 'p, 'b1, 'b2) H where
  Hmap f g  $\equiv$  Gmap (F1map f g) (F2map f g) (F3map f g)

abbreviation Hset1 :: ('p1, 'p2, 'p3, 'p, 'a1, 'a2) H  $\Rightarrow$  'a1 set where
  Hset1  $\equiv$  Union o Gset o Gmap F1set1 F2set1 F3set1

abbreviation Hset2 :: ('p1, 'p2, 'p3, 'p, 'a1, 'a2) H  $\Rightarrow$  'a2 set where
  Hset2  $\equiv$  Union o Gset o Gmap F1set2 F2set2 F3set2

lemma Hset1_alt:
  Hset1 = Union o BNF_Def.collect {image F1set1 o Gset1, image F2set1 o Gset2, image F3set1 o Gset3}
  by (tactic  $\ll$  BNF_Comp_Tactics.mk_comp_set_alt_tac  $\@$ {context}  $\@$ {thm G.collect_set_map}  $\gg$ )

lemma Hset2_alt:
  Hset2 = Union o BNF_Def.collect {image F1set2 o Gset1, image F2set2 o Gset2, image F3set2 o Gset3}
  by (tactic  $\ll$  BNF_Comp_Tactics.mk_comp_set_alt_tac  $\@$ {context}  $\@$ {thm G.collect_set_map}  $\gg$ )

```

abbreviation *Hbd* where

$$Hbd \equiv (bd_F1 + c\ bd_F2 + c\ bd_F3) *c\ bd_G$$

theorem *Hmap_id*: $Hmap\ id\ id = id$

unfolding $G.map_id0\ F1.map_id0\ F2.map_id0\ F3.map_id0\ ..$

theorem *Hmap_comp*: $Hmap\ (f1\ o\ g1)\ (f2\ o\ g2) = Hmap\ f1\ f2\ o\ Hmap\ g1\ g2$

unfolding $G.map_comp0\ F1.map_comp0\ F2.map_comp0\ F3.map_comp0\ ..$

theorem *Hmap_cong*: $\llbracket \bigwedge z. z \in Hset1\ x \implies f1\ z = g1\ z; \bigwedge z. z \in Hset2\ x \implies f2\ z = g2\ z \rrbracket \implies Hmap\ f1\ f2\ x = Hmap\ g1\ g2\ x$

by (tactic $\ll BNF_Comp_Tactics.mk_comp_map_cong0_tac\ @\{context\}$
 $\llbracket @\{thms\ Hset1_alt\ Hset2_alt\} @\{thm\ G.map_cong0\} @\{thms\ F1.map_cong0\ F2.map_cong0\ F3.map_cong0\}$
 \rrbracket)

theorem *Hset1_natural*: $Hset1\ o\ Hmap\ f1\ f2 = image\ f1\ o\ Hset1$

by (tactic $\ll BNF_Comp_Tactics.mk_comp_set_map0_tac\ @\{context\}\ @\{thm\ refl\}\ @\{thm\ G.map_comp0\}$
 $@\{thm\ G.map_cong0\}$
 $@\{thm\ G.collect_set_map\}\ @\{thms\ F1.set_map0(1)\ F2.set_map0(1)\ F3.set_map0(1)\}$)

theorem *Hset2_natural*: $Hset2\ o\ Hmap\ f1\ f2 = image\ f2\ o\ Hset2$

by (tactic $\ll BNF_Comp_Tactics.mk_comp_set_map0_tac\ @\{context\}\ @\{thm\ refl\}\ @\{thm\ G.map_comp0\}$
 $@\{thm\ G.map_cong0\}$
 $@\{thm\ G.collect_set_map\}\ @\{thms\ F1.set_map0(2)\ F2.set_map0(2)\ F3.set_map0(2)\}$)

theorem *Hbd_card_order*: $card_order\ Hbd$

by (tactic $\ll BNF_Comp_Tactics.mk_comp_bd_card_order_tac\ @\{context\}$
 $@\{thms\ F1.bd_card_order\ F2.bd_card_order\ F3.bd_card_order\}\ @\{thm\ G.bd_card_order\}$)

theorem *Hbd_cinfinite*: $cinfinite\ Hbd$

by (tactic $\ll BNF_Comp_Tactics.mk_comp_bd_cinfinite_tac\ @\{context\}$
 $@\{thm\ F1.bd_cinfinite\}\ @\{thm\ G.bd_cinfinite\}$)

theorem *Hset1_bd*: $|Hset1\ (x :: ('p1, 'p2, 'p3, 'p, 'a1, 'a2)\ H)| \leq o$

($Hbd :: ('p1, 'p2, 'p3, 'p)\ Hbd_type\ rel$)

by (tactic $\ll BNF_Comp_Tactics.mk_comp_set_bd_tac\ @\{context\}\ @\{thm\ refl\}\ NONE\ @\{thm\ Hset1_alt\}$
 $@\{thms\ comp_single_set_bd[OF\ F1.bd_Card_order\ F1.set_bd(1)\ G.set_bd(1)]$
 $comp_single_set_bd[OF\ F2.bd_Card_order\ F2.set_bd(1)\ G.set_bd(2)]$
 $comp_single_set_bd[OF\ F3.bd_Card_order\ F3.set_bd(1)\ G.set_bd(3)]\}$)

theorem *Hset2_bd*: $|Hset2\ (x :: ('p1, 'p2, 'p3, 'p, 'a1, 'a2)\ H)| \leq o$

($Hbd :: ('p1, 'p2, 'p3, 'p)\ Hbd_type\ rel$)

by (tactic $\ll BNF_Comp_Tactics.mk_comp_set_bd_tac\ @\{context\}\ @\{thm\ refl\}\ NONE\ @\{thm\ Hset2_alt\}$
 $@\{thms\ comp_single_set_bd[OF\ F1.bd_Card_order\ F1.set_bd(2)\ G.set_bd(1)]$
 $comp_single_set_bd[OF\ F2.bd_Card_order\ F2.set_bd(2)\ G.set_bd(2)]$
 $comp_single_set_bd[OF\ F3.bd_Card_order\ F3.set_bd(2)\ G.set_bd(3)]\}$)

abbreviation *Hin* where $Hin\ A1\ A2 \equiv \{x. Hset1\ x \subseteq A1 \wedge Hset2\ x \subseteq A2\}$

lemma *Hin_alt*: $Hin\ A1\ A2 = Gin\ (F1in\ A1\ A2)\ (F2in\ A1\ A2)\ (F3in\ A1\ A2)$

by (tactic $\ll BNF_Comp_Tactics.mk_comp_in_alt_tac\ @\{context\}\ @\{thms\ Hset1_alt\ Hset2_alt\}$)

definition *Hwit1* where $Hwit1\ b\ c = wit1_G\ wit_F1\ (wit_F3\ b\ c)$

definition *Hwit21* where $Hwit21\ b\ c = wit2_G\ (wit1_F2\ b)\ (wit_F3\ b\ c)$

definition *Hwit22* where $Hwit22\ b\ c = wit2_G\ (wit2_F2\ c)\ (wit_F3\ b\ c)$

lemma *Hwit1*:

$\bigwedge x. x \in Hset1\ (Hwit1\ b\ c) \implies x = b$

$\bigwedge x. x \in Hset2\ (Hwit1\ b\ c) \implies x = c$

unfolding *Hwit1_def*

by (tactic $\ll BNF_Comp_Tactics.mk_comp_wit_tac\ @\{context\}\ \llbracket @\{thms\ G.wit1\ G.wit2\}$
 $@\{thm\ G.collect_set_map\}\ @\{thms\ F1.wit\ F2.wit1\ F2.wit2\ F3.wit\}$)

lemma *Hwit21*:

$\bigwedge x. x \in \text{Hset1 } (\text{Hwit21 } b \ c) \implies x = b$
 $\bigwedge x. x \in \text{Hset2 } (\text{Hwit21 } b \ c) \implies x = c$

unfolding *Hwit21_def*

by (*tactic* $\ll \text{BNF_Comp_Tactics.mk_comp_wit_tac } @\{\text{context}\} \ \square \ @\{\text{thms } G.\text{wit1 } G.\text{wit2}\} @\{\text{thm } G.\text{collect_set_map}\} @\{\text{thms } F1.\text{wit } F2.\text{wit1 } F2.\text{wit2 } F3.\text{wit}\} \gg$)

lemma *Hwit22*:

$\bigwedge x. x \in \text{Hset1 } (\text{Hwit22 } b \ c) \implies x = b$
 $\bigwedge x. x \in \text{Hset2 } (\text{Hwit22 } b \ c) \implies x = c$

unfolding *Hwit22_def*

by (*tactic* $\ll \text{BNF_Comp_Tactics.mk_comp_wit_tac } @\{\text{context}\} \ \square \ @\{\text{thms } G.\text{wit1 } G.\text{wit2}\} @\{\text{thm } G.\text{collect_set_map}\} @\{\text{thms } F1.\text{wit } F2.\text{wit1 } F2.\text{wit2 } F3.\text{wit}\} \gg$)

lemma *Grel_cong*: $[[R1 = S1; R2 = S2; R3 = S3]] \implies \text{Grel } R1 \ R2 \ R3 = \text{Grel } S1 \ S2 \ S3$

by *hypsubst* (*rule refl*)

definition *Hrel where*

Hrel $R1 \ R2 = (\text{BNF_Def.Grp } (\text{Hin } (\text{Collect } (\text{case_prod } R1)) (\text{Collect } (\text{case_prod } R2)))) (\text{Hmap fst fst})^{--1} \text{OO}$

$(\text{BNF_Def.Grp } (\text{Hin } (\text{Collect } (\text{case_prod } R1)) (\text{Collect } (\text{case_prod } R2)))) (\text{Hmap snd snd})$

lemmas *Hrel_unfold* = *trans*[*OF Hrel_def trans*[*OF OO_Grp_cong*[*OF Hin_alt*

trans[*OF arg_cong2*[*of _ _ _ relcompp*, *OF trans*[*OF arg_cong*[*of _ _ conversep*, *OF sym*[*OF G.rel_Grp*]]

G.rel_conversep[*symmetric*]] *sym*[*OF G.rel_Grp*]] *trans*[*OF G.rel_compp*[*symmetric*] *Grel_cong*[*OF sym*[*OF F1.rel_compp_Grp*] *sym*[*OF F2.rel_compp_Grp*] *sym*[*OF F3.rel_compp_Grp*]]]]]

bnf *H*: (*'p1*, *'p2*, *'p3*, *'p*, *'a1*, *'a2*) *H*

map: *Hmap*

sets: *Hset1 Hset2*

bd: *Hbd* :: (*'p1*, *'p2*, *'p3*, *'p*) *Hbd_type rel*

rel: *Hrel*

apply –

apply (*rule Hmap_id*)

apply (*rule Hmap_comp*)

apply (*erule Hmap_cong*) **apply** *assumption*

apply (*rule Hset1_natural*)

apply (*rule Hset2_natural*)

apply (*rule Hbd_card_order*)

apply (*rule Hbd_cinfinite*)

apply (*rule Hset1_bd*)

apply (*rule Hset2_bd*)

apply (*unfold Hrel_unfold* *G.rel_compp*[*symmetric*] *F1.rel_compp*[*symmetric*] *F2.rel_compp*[*symmetric*] *F3.rel_compp*[*symmetric*] *eq_OO*] [*1*] **apply** (*rule order_refl*)

apply (*rule Hrel_def*[*unfolded OO_Grp_alt mem_Collect_eq*])

done

4 Removing Live Variables

declare [*bnf_internals*]

bnf-axiomatization (*dead 'p*, *Fset1*: *'a1*, *Fset2*: *'a2*, *Fset3*: *'a3*) *F* **for** *map*: *Fmap* *rel*: *Frel*

abbreviation *F1map* :: (*'a2* \implies *'b2*) \implies (*'a3* \implies *'b3*) \implies (*'p*, *'a1*, *'a2*, *'a3*) *F* \implies (*'p*, *'a1*, *'b2*, *'b3*) *F* **where**

F1map \equiv *Fmap id*

abbreviation *F2map* :: (*'a3* \implies *'b3*) \implies (*'p*, *'a1*, *'a2*, *'a3*) *F* \implies (*'p*, *'a1*, *'a2*, *'b3*) *F* **where**

F2map \equiv *Fmap id id*

abbreviation $F1set1 \equiv Fset2$
abbreviation $F1set2 \equiv Fset3$
abbreviation $F2set \equiv Fset3$

theorem $F1map_id$: $F1map\ id\ id = id$
by (rule $F.map_id0$)

theorem $F2map_id$: $F2map\ id = id$
by (rule $F.map_id0$)

theorem $F1map_comp$: $F1map\ (f1\ o\ g1)\ (f2\ o\ g2) = F1map\ f1\ f2\ o\ F1map\ g1\ g2$
by (unfold $F.map_comp0[symmetric]$ o_id) (rule $refl$)

theorem $F2map_comp$: $F2map\ (f\ o\ g) = F2map\ f\ o\ F2map\ g$
by (unfold $F.map_comp0[symmetric]$ o_id) (rule $refl$)

theorem $F1map_cong$: $[\bigwedge z. z \in F1set1\ x \implies f1\ z = g1\ z; \bigwedge z. z \in F1set2\ x \implies f2\ z = g2\ z] \implies F1map\ f1\ f2\ x = F1map\ g1\ g2\ x$
apply (rule $F.map_cong0$)
apply (rule $refl$)
apply $assumption$
apply $assumption$
done

theorem $F2map_cong$: $[\bigwedge z. z \in F2set\ x \implies f\ z = g\ z] \implies F2map\ f\ x = F2map\ g\ x$
apply (rule $F.map_cong0$)
apply (rule $refl$)
apply (rule $refl$)
apply $assumption$
done

theorem $F1set1_natural$: $F1set1\ o\ F1map\ f1\ f2 = image\ f1\ o\ F1set1$
by (rule $F.set_map0(2)$)

theorem $F1set2_natural$: $F1set2\ o\ F1map\ f1\ f2 = image\ f2\ o\ F1set2$
by (rule $F.set_map0(3)$)

theorem $F2set_natural$: $F2set\ o\ F2map\ f = image\ f\ o\ F2set$
by (rule $F.set_map0(3)$)

abbreviation $Fin :: 'a1\ set \Rightarrow 'a2\ set \Rightarrow 'a3\ set \Rightarrow ((p, 'a1, 'a2, 'a3)\ F)\ set$ **where**
 $Fin\ A1\ A2\ A3 \equiv \{x. Fset1\ x \subseteq A1 \wedge Fset2\ x \subseteq A2 \wedge Fset3\ x \subseteq A3\}$

abbreviation $F1in :: 'a2\ set \Rightarrow 'a3\ set \Rightarrow ((p, 'a1, 'a2, 'a3)\ F)\ set$ **where**
 $F1in\ A1\ A2 \equiv \{x. F1set1\ x \subseteq A1 \wedge F1set2\ x \subseteq A2\}$

lemma $F1in_alt$: $F1in\ A2\ A3 = Fin\ UNIV\ A2\ A3$
by (tactic $\langle BNF_Comp_Tactics.kill_in_alt_tac\ @\{context\} \rangle$)

abbreviation $F2in :: 'a3\ set \Rightarrow ((p, 'a1, 'a2, 'a3)\ F)\ set$ **where**
 $F2in\ A \equiv \{x. F2set\ x \subseteq A\}$

lemma $F2in_alt$: $F2in\ A3 = Fin\ UNIV\ UNIV\ A3$
by (tactic $\langle BNF_Comp_Tactics.kill_in_alt_tac\ @\{context\} \rangle$)

lemma $Frel_cong$: $[\mathbb{R}1 = \mathbb{S}1; \mathbb{R}2 = \mathbb{S}2; \mathbb{R}3 = \mathbb{S}3] \implies Frel\ \mathbb{R}1\ \mathbb{R}2\ \mathbb{R}3 = Frel\ \mathbb{S}1\ \mathbb{S}2\ \mathbb{S}3$
by $hypsubst$ (rule $refl$)

definition $F1rel$ **where**
 $F1rel\ \mathbb{R}1\ \mathbb{R}2 = (BNF_Def.Grp\ (F1in\ (Collect\ (case_prod\ \mathbb{R}1))\ (Collect\ (case_prod\ \mathbb{R}2))))\ (F1map\ fst\ fst))\ ^{--1}$
 OO

(BNF_Def.Grp (F1in (Collect (case_prod R1)) (Collect (case_prod R2))) (F1map snd snd))

lemmas F1rel_unfold = trans[OF F1rel_def trans[OF OO_Grp_cong[OF F1in_alt]
trans[OF arg_cong2[of _ _ _ relcompp, OF trans[OF arg_cong[of _ _ conversep, OF sym[OF F.rel_Grp]]
F.rel_conversep[symmetric]] sym[OF F.rel_Grp]]
trans[OF F.rel_compp[symmetric] Frel_cong[OF trans[OF Grp_UNIV_id[OF refl] eq_alt[symmetric]]
Grp_fst_snd Grp_fst_snd]]]]]]

definition F2rel where

F2rel R1 = (BNF_Def.Grp (F2in (Collect (case_prod R1))) (F2map fst)) ^--1 OO
(BNF_Def.Grp (F2in (Collect (case_prod R1))) (F2map snd))

lemmas F2rel_unfold = trans[OF F2rel_def trans[OF OO_Grp_cong[OF F2in_alt]
trans[OF arg_cong2[of _ _ _ relcompp, OF trans[OF arg_cong[of _ _ conversep, OF sym[OF F.rel_Grp]]
F.rel_conversep[symmetric]] sym[OF F.rel_Grp]]
trans[OF F.rel_compp[symmetric] Frel_cong[OF trans[OF Grp_UNIV_id[OF refl] eq_alt[symmetric]]
trans[OF Grp_UNIV_id[OF refl] eq_alt[symmetric]] Grp_fst_snd]]]]]]

bnf F1: ('p, 'a1, 'a2, 'a3) F

map: F1map

sets: F1set1 F1set2

bd: bd_F :: ('p bd_type_F) rel

rel: F1rel

apply –
apply (rule F1map_id)
apply (rule F1map_comp)
apply (erule F1map_cong) **apply** assumption
apply (rule F1set1_natural)
apply (rule F1set2_natural)
apply (rule F.bd_card_order)
apply (rule F.bd_cinfinite)
apply (rule F.set_bd(2))
apply (rule F.set_bd(3))
apply (unfold F1rel_unfold F.rel_compp[symmetric] eq_OO) [1] **apply** (rule order_refl)
apply (rule F1rel_def[unfolded OO_Grp_alt mem_Collect_eq])
done

bnf F2: ('p, 'a1, 'a2, 'a3) F

map: F2map

sets: F2set

bd: bd_F :: ('p bd_type_F) rel

rel: F2rel

apply –
apply (rule F2map_id)
apply (rule F2map_comp)
apply (erule F2map_cong)
apply (rule F2set_natural)
apply (rule F.bd_card_order)
apply (rule F.bd_cinfinite)
apply (rule F.set_bd(3))
apply (unfold F2rel_unfold F.rel_compp[symmetric] eq_OO) [1] **apply** (rule order_refl)
apply (rule F2rel_def[unfolded OO_Grp_alt mem_Collect_eq])
done

5 Adding New Live Variables

declare [[bnf_internals]]

bnf-axiomatization (dead 'p, Fset1: 'a1, Fset2: 'a2) F

[wits: 'a1 \Rightarrow 'a2 \Rightarrow ('p, 'a1, 'a2) F]

for map: Fmap rel: Frel

type-synonym ('p, 'a1, 'a2, 'a3, 'a4) F' = ('p, 'a3, 'a4) F

abbreviation $F'_{map} :: ('a1 \Rightarrow 'b1) \Rightarrow ('a2 \Rightarrow 'b2) \Rightarrow ('a3 \Rightarrow 'b3) \Rightarrow ('a4 \Rightarrow 'b4) \Rightarrow ('p, 'a1, 'a2, 'a3, 'a4) F' \Rightarrow ('p, 'b1, 'b2, 'b3, 'b4) F'$ **where**
 $F'_{map} f1 f2 f3 f4 \equiv F_{map} f3 f4$

abbreviation $F'_{set1} :: ('p, 'a1, 'a2, 'a3, 'a4) F' \Rightarrow 'a1$ **set where**
 $F'_{set1} \equiv \lambda_ . \{\}$

abbreviation $F'_{set2} :: ('p, 'a1, 'a2, 'a3, 'a4) F' \Rightarrow 'a2$ **set where**
 $F'_{set2} \equiv \lambda_ . \{\}$

abbreviation $F'_{set3} :: ('p, 'a1, 'a2, 'a3, 'a4) F' \Rightarrow 'a3$ **set where**
 $F'_{set3} \equiv F_{set1}$

abbreviation $F'_{set4} :: ('p, 'a1, 'a2, 'a3, 'a4) F' \Rightarrow 'a4$ **set where**
 $F'_{set4} \equiv F_{set2}$

abbreviation F'_{bd} **where**
 $F'_{bd} \equiv bd_F$

theorem F'_{map_id} : $F'_{map} id id id id = id$
by (rule F_{map_id0})

theorem F'_{map_comp} :
 $F'_{map} (f1 o g1) (f2 o g2) (f3 o g3) (f4 o g4) = F'_{map} f1 f2 f3 f4 o F'_{map} g1 g2 g3 g4$
by (rule F_{map_comp0})

theorem F'_{map_cong} :
 $\llbracket \bigwedge z. z \in F'_{set1} x \implies f1 z = g1 z; \bigwedge z. z \in F'_{set2} x \implies f2 z = g2 z; \bigwedge z. z \in F'_{set3} x \implies f3 z = g3 z; \bigwedge z. z \in F'_{set4} x \implies f4 z = g4 z \rrbracket$
 $\implies F'_{map} f1 f2 f3 f4 x = F'_{map} g1 g2 g3 g4 x$
apply (tactic $\ll BNF_Util.rtac @\{context\} @\{thm F_{map_cong}\} 1 THEN REPEAT_DETERM_N 2 (assume_tac @\{context\} 1)\rrbracket$)
apply *assumption+*
done

theorem $F'_{set1_natural}$: $F'_{set1} o F'_{map} f1 f2 f3 f4 = image f1 o F'_{set1}$
by (tactic $\langle BNF_Comp_Tactics.empty_natural_tac @\{context\} \rangle$)

theorem $F'_{set2_natural}$: $F'_{set2} o F'_{map} f1 f2 f3 f4 = image f2 o F'_{set2}$
by (tactic $\langle BNF_Comp_Tactics.empty_natural_tac @\{context\} \rangle$)

theorem $F'_{set3_natural}$: $F'_{set3} o F'_{map} f1 f2 f3 f4 = image f3 o F'_{set3}$
by (rule $F_{set_map0}(1)$)

theorem $F'_{set4_natural}$: $F'_{set4} o F'_{map} f1 f2 f3 f4 = image f4 o F'_{set4}$
by (rule $F_{set_map0}(2)$)

theorem $F'_{bd_card_order}$: $card_order bd_F$
by (rule $F_{bd_card_order}$)

theorem $F'_{bd_cinfinte}$: $cinfinte bd_F$
by (rule $F_{bd_cinfinte}$)

theorem F'_{set1_bd} : $|F'_{set1} x| \leq o F'_{bd}$
by (tactic $\ll BNF_Comp_Tactics.mk_lift_set_bd_tac @\{context\} @\{thm F_{bd_Card_order}\} \rrbracket$)

theorem F'_{set2_bd} : $|F'_{set2} x| \leq o F'_{bd}$
by (tactic $\ll BNF_Comp_Tactics.mk_lift_set_bd_tac @\{context\} @\{thm F_{bd_Card_order}\} \rrbracket$)

theorem F'_{set3_bd} : $|F'_{set3} (x :: ('c, 'a, 'd) F)| \leq o (F'_{bd} :: 'c bd_type_F rel)$
by (rule $F_{set_bd}(1)$)

theorem $F'set4_bd$: $|F'set4 (x :: ('c, 'a, 'd) F)| \leq_o (F'bd :: 'c\ bd_type_F\ rel)$
by (rule $F.set_bd(2)$)

abbreviation $F'in :: 'a1\ set \Rightarrow 'a2\ set \Rightarrow 'a3\ set \Rightarrow 'a4\ set \Rightarrow ((p, 'a1, 'a2, 'a3, 'a4) F)$ **set** **where**
 $F'in\ A1\ A2\ A3\ A4 \equiv \{x. F'set1\ x \subseteq A1 \wedge F'set2\ x \subseteq A2 \wedge F'set3\ x \subseteq A3 \wedge F'set4\ x \subseteq A4\}$

definition $F'rel$ **where**

$$F'rel\ R1\ R2\ R3\ R4 = (BNF_Def.Grp (F'in (Collect (case_prod\ R1))) (Collect (case_prod\ R2))) (Collect (case_prod\ R3)) (Collect (case_prod\ R4)) (F'map\ fst\ fst\ fst\ fst) \hat{---} 1\ OO$$

$$(BNF_Def.Grp (F'in (Collect (case_prod\ R1))) (Collect (case_prod\ R2))) (Collect (case_prod\ R3)) (Collect (case_prod\ R4)) (F'map\ snd\ snd\ snd\ snd))$$

lemmas $F'rel_unfold = trans[OF\ F'rel_def[unfolded\ eqTrueI[OF\ empty_subsetI]\ simp_thms(22)]$
 $trans[OF\ OO_Grp_cong[OF\ refl]\ sym[OF\ F.rel_compp_Grp]]]$

bnf F' : $(p, 'a1, 'a2, 'a3, 'a4) F'$
 map : $F'map$
 $sets$: $F'set1\ F'set2\ F'set3\ F'set4$
 bd : $F'bd :: 'p\ bd_type_F\ rel$
 $wits$: wit_F
 rel : $F'rel$
 $plugins\ del$: $lifting\ transfer$
apply –
apply (rule $F'map_id$)
apply (rule $F'map_comp$)
apply (erule $F'map_cong$) **apply** $assumption+$
apply (rule $F'set1_natural$)
apply (rule $F'set2_natural$)
apply (rule $F'set3_natural$)
apply (rule $F'set4_natural$)
apply (rule $F'bd_card_order$)
apply (rule $F'bd_cinfinte$)
apply (rule $F'set1_bd$)
apply (rule $F'set2_bd$)
apply (rule $F'set3_bd$)
apply (rule $F'set4_bd$)
apply (unfold $F'rel_unfold\ F.rel_compp[symmetric]\ eq_OO$) [1] **apply** (rule $order_refl$)
apply (rule $F'rel_def[unfolded\ OO_Grp_alt\ mem_Collect_eq]$)
apply (erule $F.wit_emptyE$)
done

6 Changing the Order of Live Variables

declare $[[bnf_internals]]$

bnf-axiomatization ($dead\ 'p, Fset1: 'a1, Fset2: 'a2, Fset3: 'a3$) F **for** map : $Fmap$ rel : $Frel$

type-synonym $(p, 'a1, 'a2, 'a3) F' = (p, 'a3, 'a1, 'a2) F$

abbreviation $Fin :: 'a1\ set \Rightarrow 'a2\ set \Rightarrow 'a3\ set \Rightarrow ((p, 'a1, 'a2, 'a3) F)$ **set** **where**

$$Fin\ A1\ A2\ A3 \equiv \{x. Fset1\ x \subseteq A1 \wedge Fset2\ x \subseteq A2 \wedge Fset3\ x \subseteq A3\}$$

abbreviation $F'map :: ('a1 \Rightarrow 'b1) \Rightarrow ('a2 \Rightarrow 'b2) \Rightarrow ('a3 \Rightarrow 'b3) \Rightarrow (p, 'a1, 'a2, 'a3) F' \Rightarrow (p, 'b1, 'b2, 'b3) F'$ **where**

$$F'map\ f\ g\ h \equiv Fmap\ h\ f\ g$$

abbreviation $F'set1 :: (p, 'a1, 'a2, 'a3) F' \Rightarrow 'a1\ set$ **where**

$$F'set1 \equiv Fset2$$

abbreviation $F'set2 :: (p, 'a1, 'a2, 'a3) F' \Rightarrow 'a2\ set$ **where**

$$F'set2 \equiv Fset3$$

abbreviation $F'set3 :: (p, 'a1, 'a2, 'a3) F' \Rightarrow 'a3\ set$ **where**

$F'set3 \equiv Fset1$

abbreviation $F'bd$ **where**

$F'bd \equiv bd_F$

theorem $F'map_id$: $F'map\ id\ id\ id = id$

by (rule $F.map_id0$)

theorem $F'map_comp$: $F'map\ (f1\ o\ g1)\ (f2\ o\ g2)\ (f3\ o\ g3) = F'map\ f1\ f2\ f3\ o\ F'map\ g1\ g2\ g3$

by (rule $F.map_comp0$)

theorem $F'map_cong$: $\llbracket \bigwedge z. z \in F'set1\ x \implies f1\ z = g1\ z; \bigwedge z. z \in F'set2\ x \implies f2\ z = g2\ z; \bigwedge z. z \in F'set3\ x \implies f3\ z = g3\ z \rrbracket$

$\implies F'map\ f1\ f2\ f3\ x = F'map\ g1\ g2\ g3\ x$

apply (rule $F.map_cong0$)

apply $assumption+$

done

theorem $F'set1_natural$: $F'set1\ o\ F'map\ f1\ f2\ f3 = image\ f1\ o\ F'set1$

by (rule $F.set_map0(2)$)

theorem $F'set2_natural$: $F'set2\ o\ F'map\ f1\ f2\ f3 = image\ f2\ o\ F'set2$

by (rule $F.set_map0(3)$)

theorem $F'set3_natural$: $F'set3\ o\ F'map\ f1\ f2\ f3 = image\ f3\ o\ F'set3$

by (rule $F.set_map0(1)$)

theorem $F'bd_card_order$: $card_order\ F'bd$

by (rule $F.bd_card_order$)

theorem $F'bd_cinfinte$: $cinfinte\ F'bd$

by (rule $F.bd_cinfinte$)

theorem $F'set1_bd$: $|F'set1\ (x :: ('c, 'a, 'b, 'd)\ F)| \leq o\ (F'bd :: 'c\ bd_type_F\ rel)$

by (rule $F.set_bd(2)$)

theorem $F'set2_bd$: $|F'set2\ (x :: ('c, 'a, 'b, 'd)\ F)| \leq o\ (F'bd :: 'c\ bd_type_F\ rel)$

by (rule $F.set_bd(3)$)

theorem $F'set3_bd$: $|F'set3\ (x :: ('c, 'a, 'b, 'd)\ F)| \leq o\ (F'bd :: 'c\ bd_type_F\ rel)$

by (rule $F.set_bd(1)$)

abbreviation $F'in :: 'a1\ set \Rightarrow 'a2\ set \Rightarrow 'a3\ set \Rightarrow ((p, 'a1, 'a2, 'a3)\ F)\ set$ **where**

$F'in\ A1\ A2\ A3 \equiv \{x. F'set1\ x \subseteq A1 \wedge F'set2\ x \subseteq A2 \wedge F'set3\ x \subseteq A3\}$

lemma $F'in_alt$: $F'in\ A1\ A2\ A3 = Fin\ A3\ A1\ A2$

apply (rule $Collect_cong$)

by (tactic $\ll BNF_Tactics.mk_rotate_eq_tac\ @\{context\}$

$(BNF_Util.rtac\ @\{context\}\ @\{thm\ refl\})\ @\{thm\ trans\}\ @\{thm\ conj_assoc\}\ @\{thm\ conj_commute\}\ @\{thm\ conj_cong\}$

$[1, 2, 3]\ [3, 1, 2]\ 1\ \rangle\rangle$)

definition $F'rel$ **where**

$F'rel\ R1\ R2\ R3 = (BNF_Def.Grp\ (F'in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))))\ (F'map\ fst\ fst\ fst)\ ^{-1}\ OO$

$(BNF_Def.Grp\ (F'in\ (Collect\ (case_prod\ R1))\ (Collect\ (case_prod\ R2))\ (Collect\ (case_prod\ R3))))\ (F'map\ snd\ snd\ snd)$

lemmas $F'rel_unfold = trans[OF\ F'rel_def\ trans[OF\ OO_Grp_cong[OF\ F'in_alt]\ sym[OF\ F.rel_compp_Grp]]]$

bnf F' : $(p, 'a1, 'a2, 'a3)\ F'$

$map: F'map$


```

sets: F'set1 F'set2 F'set3
bd: F'bd :: 'p bd_type_F rel
rel: F'rel
  apply -
  apply (rule F'map_id)
  apply (rule F'map_comp)
  apply (erule F'map_cong) apply assumption+
  apply (rule F'set1_natural)
  apply (rule F'set2_natural)
  apply (rule F'set3_natural)
  apply (rule F'bd_card_order)
  apply (rule F'bd_cinfinite)
  apply (rule F'set1_bd)
  apply (rule F'set2_bd)
  apply (rule F'set3_bd)
  apply (unfold F'rel_unfold F.rel_compp[symmetric] eq_OO) [1] apply (rule order_refl)
  apply (rule F'rel_def[unfolded OO_Grp_alt mem_Collect_eq])
done

```

7 Mutual View on Nested Datatypes

notation *BNF_Def.convolve* ($\langle _ , _ \rangle$)

declare *[[bnf_internals]]*

declare *[[typedef_overloaded]]*

bnf-axiomatization ('a, 'b) F0 [wits: 'a \Rightarrow ('a, 'b) F0]

bnf-axiomatization ('a, 'b) G0 [wits: 'a \Rightarrow ('a, 'b) G0]

7.1 Nested Definition

datatype 'a F = CF ('a, 'a F) F0

datatype 'a G = CG ('a, ('a G) F) G0

type-synonym ('b, 'c) F_pre_F = ('c, 'b) F0

type-synonym ('c, 'a) G_pre_G = ('a, 'c F) G0

term *ctor_fold_F* :: (('b, 'c) F_pre_F \Rightarrow 'b) \Rightarrow 'c F \Rightarrow 'b

term *ctor_fold_G* :: (('c, 'a) G_pre_G \Rightarrow 'c) \Rightarrow 'a G \Rightarrow 'c

term *ctor_rec_F* :: (('c F \times 'b, 'c) F_pre_F \Rightarrow 'b) \Rightarrow 'c F \Rightarrow 'b

term *ctor_rec_G* :: (('a G \times 'c, 'a) G_pre_G \Rightarrow 'c) \Rightarrow 'a G \Rightarrow 'c

thm *F.ctor_rel_induct*

thm *G.ctor_rel_induct[unfolded rel_pre_G_def id_apply]*

7.2 Isomorphic Mutual Definition

datatype 'a G_M = CG ('a, 'a GF_M) G0

and 'a GF_M = CF ('a G_M, 'a GF_M) F0

type-synonym ('b, 'c) GF_M_pre_GF_M = ('c, 'b) F0

type-synonym ('c, 'a) G_M_pre_G_M = ('a, 'c) G0

term *ctor_fold_G_M* :: (('c, 'a) G_M_pre_G_M \Rightarrow 'b) \Rightarrow (('c, 'b) GF_M_pre_GF_M \Rightarrow 'c) \Rightarrow 'a G_M \Rightarrow 'b

term *ctor_fold_GF_M* :: (('c, 'a) G_M_pre_G_M \Rightarrow 'b) \Rightarrow (('c, 'b) GF_M_pre_GF_M \Rightarrow 'c) \Rightarrow 'a GF_M \Rightarrow 'c

term *ctor_rec_G_M* :: (('a GF_M \times 'c, 'a) G_M_pre_G_M \Rightarrow 'b) \Rightarrow (('a GF_M \times 'c, 'a G_M \times 'b) GF_M_pre_GF_M \Rightarrow 'c) \Rightarrow 'a G_M \Rightarrow 'b

term *ctor_rec_GF_M* :: (('a GF_M \times 'c, 'a) G_M_pre_G_M \Rightarrow 'b) \Rightarrow (('a GF_M \times 'c, 'a G_M \times 'b) GF_M_pre_GF_M \Rightarrow 'c) \Rightarrow 'a GF_M \Rightarrow 'c

thm *G_M_GF_M.ctor_rel_induct[unfolded rel_pre_G_M_def rel_pre_GF_M_def]*

7.3 Mutualization

7.3.1 Iterators

definition $n2m_ctor_fold_G :: (('c, 'a) G_M_pre_G_M \Rightarrow 'b) \Rightarrow (('c, 'b) GF_M_pre_GF_M \Rightarrow 'c) \Rightarrow 'a G \Rightarrow 'b$
where $n2m_ctor_fold_G s1 s2 = ctor_fold_G (s1 o$
 $map_pre_G_M id (id :: unit \Rightarrow unit) (ctor_fold_F (s2 o BNF_Composition.id_bnf o BNF_Composition.id_bnf))$
 $o BNF_Composition.id_bnf o BNF_Composition.id_bnf)$

definition $n2m_ctor_fold_G_F :: (('c, 'a) G_M_pre_G_M \Rightarrow 'b) \Rightarrow (('c, 'b) GF_M_pre_GF_M \Rightarrow 'c) \Rightarrow 'a G F$
 $\Rightarrow 'c$
where $n2m_ctor_fold_G_F s1 s2 = ctor_fold_F (s2 o map_pre_GF_M (id :: unit \Rightarrow unit) (n2m_ctor_fold_G$
 $s1 s2) id o BNF_Composition.id_bnf o BNF_Composition.id_bnf)$

lemma $G_ctor_o_fold: ctor_fold_G s o ctor_G = s o map_pre_G id (ctor_fold_G s)$
unfolding $fun_eq_iff o_apply G.ctor_fold$ **by** $simp$

lemma $F_ctor_o_fold: ctor_fold_F s o ctor_F = s o map_pre_F id (ctor_fold_F s)$
unfolding $fun_eq_iff o_apply F.ctor_fold$ **by** $simp$

lemma $G_ctor_o_rec: ctor_rec_G s o ctor_G = s o map_pre_G id (BNF_Def.convolve id (ctor_rec_G s))$
unfolding $fun_eq_iff o_apply G.ctor_rec$ **by** $simp$

lemma $F_ctor_o_rec: ctor_rec_F s o ctor_F = s o map_pre_F id (BNF_Def.convolve id (ctor_rec_F s))$
unfolding $fun_eq_iff o_apply F.ctor_rec$ **by** $simp$

lemma $n2m_ctor_fold_G:$
 $n2m_ctor_fold_G s1 s2 o ctor_G = s1 o map_pre_G_M id id (n2m_ctor_fold_G_F s1 s2) o BNF_Composition.id_bnf$
 $o BNF_Composition.id_bnf$

unfolding $n2m_ctor_fold_G_def n2m_ctor_fold_G_F_def$
 $map_pre_G_def map_pre_F_def map_pre_G_M_def map_pre_GF_M_def$
 $G_ctor_o_fold id apply comp id id comp comp assoc$
 $rewriteL_comp_comp[OF type_copy_map_comp0_undo[OF BNF_Composition.type_definition_id_bnf_UNIV$
 $BNF_Composition.type_definition_id_bnf_UNIV BNF_Composition.type_definition_id_bnf_UNIV pre_G_M.map_comp0[unfold$
 $map_pre_G_M_def]]]$
 $F.ctor_fold_o_map$
 $rewriteL_comp_comp[OF type_copy_Rep_o_Abs[OF BNF_Composition.type_definition_id_bnf_UNIV]] ..$

lemma $n2m_ctor_fold_G_F:$
 $n2m_ctor_fold_G_F s1 s2 o ctor_F = s2 o map_pre_GF_M id (n2m_ctor_fold_G s1 s2) (n2m_ctor_fold_G_F$
 $s1 s2) o BNF_Composition.id_bnf o BNF_Composition.id_bnf$

unfolding $n2m_ctor_fold_G_F_def map_pre_F_def map_pre_G_M_def map_pre_GF_M_def$
 $F_ctor_o_fold id apply comp id id comp comp assoc$
 $rewriteL_comp_comp[OF F0.map_comp0[symmetric]]$
 $rewriteL_comp_comp[OF type_copy_Rep_o_Abs[OF BNF_Composition.type_definition_id_bnf_UNIV]] ..$

7.3.2 Recursors

definition $n2m_ctor_rec_G ::$
 $(('a G F \times 'c, 'a) G_M_pre_G_M \Rightarrow 'b) \Rightarrow (('a G F \times 'c, 'a G \times 'b) GF_M_pre_GF_M \Rightarrow 'c) \Rightarrow 'a G \Rightarrow 'b$
where $n2m_ctor_rec_G s1 s2 =$
 $ctor_rec_G (s1 o$
 $map_pre_G_M id (id :: unit \Rightarrow unit)$
 $(BNF_Def.convolve (map_F fst) (ctor_rec_F (s2 o map_pre_GF_M (id :: unit \Rightarrow unit) id (map_prod (map_F$
 $fst) id) o BNF_Composition.id_bnf o BNF_Composition.id_bnf))) o$
 $BNF_Composition.id_bnf o BNF_Composition.id_bnf)$

definition $n2m_ctor_rec_G_F ::$
 $(('a G F \times 'c, 'a) G_M_pre_G_M \Rightarrow 'b) \Rightarrow (('a G F \times 'c, 'a G \times 'b) GF_M_pre_GF_M \Rightarrow 'c) \Rightarrow 'a G F \Rightarrow 'c$
where $n2m_ctor_rec_G_F s1 s2 = ctor_rec_F (s2 o map_pre_GF_M (id :: unit \Rightarrow unit) (BNF_Def.convolve id$
 $(n2m_ctor_rec_G s1 s2)) id o BNF_Composition.id_bnf o BNF_Composition.id_bnf)$

lemma $n2m_ctor_rec_G:$
 $n2m_ctor_rec_G s1 s2 o ctor_G = s1 o map_pre_G_M id id (BNF_Def.convolve id (n2m_ctor_rec_G_F s1 s2))$
 $o BNF_Composition.id_bnf o BNF_Composition.id_bnf$

unfolding $n2m_ctor_rec_G_def n2m_ctor_rec_G_F_def$
 $map_pre_G_def map_pre_F_def map_pre_G_M_def map_pre_GF_M_def$
 $G_ctor_o_rec$

```

id_apply comp_id id_comp comp_assoc map_prod.comp map_prod.id
fst_convolver map_prod_o_convolver convolver_o
rewriteL_comp_comp[OF G0.map_comp0[symmetric]]
rewriteL_comp_comp[OF F0.map_comp0[symmetric]]
F.map_comp0[symmetric] F.map_id0
F.ctor_rec_o_map
rewriteL_comp_comp[OF type_copy_Rep_o_Abs[OF BNF_Composition.type_definition_id_bnf_UNIV]] ..

```

lemma *n2m_ctor_rec_G_F*:
n2m_ctor_rec_G_F s1 s2 o ctor_F = s2 o map_pre_GF_M id (BNF_Def.convolver id (n2m_ctor_rec_G s1 s2))
(BNF_Def.convolver id (n2m_ctor_rec_G_F s1 s2)) o BNF_Composition.id_bnf o BNF_Composition.id_bnf
unfolding *n2m_ctor_rec_G_F_def map_pre_F_def map_pre_G_M_def map_pre_GF_M_def*
F_ctor_o_rec_id_apply_comp_id_id_comp_comp_assoc
rewriteL_comp_comp[OF F0.map_comp0[symmetric]]
rewriteL_comp_comp[OF type_copy_Rep_o_Abs[OF BNF_Composition.type_definition_id_bnf_UNIV]] ..

7.3.3 Induction

lemma *n2m_rel_induct_G_G_F*:
assumes *IH1: $\forall x y. \text{BNF_Def.vimage2p} (\text{BNF_Composition.id_bnf} \circ \text{BNF_Composition.id_bnf}) (\text{BNF_Composition.id_bnf} \circ \text{BNF_Composition.id_bnf}) (\text{rel_pre_G}_M P R S) x y \longrightarrow R (\text{ctor_G } x) (\text{ctor_G } y)$*
and *IH2: $\forall x y. \text{BNF_Def.vimage2p} (\text{BNF_Composition.id_bnf} \circ \text{BNF_Composition.id_bnf}) (\text{BNF_Composition.id_bnf} \circ \text{BNF_Composition.id_bnf}) (\text{rel_pre_GF}_M P R S) x y \longrightarrow S (\text{ctor_F } x) (\text{ctor_F } y)$*
shows *rel_G P \leq R \wedge rel_F (rel_G P) \leq S*
apply (*rule context_conjI*)
apply (*rule G.ctor_rel_induct[unfolding rel_pre_G_def id_apply vimage2p_def o_apply]*)
apply (*erule mp[OF spec2[OF IH1], OF vimage2p_mono[OF _ pre_G_M.rel_mono], unfolding vimage2p_def o_apply rel_pre_G_M_def type_definition.Abs_inverse[OF BNF_Composition.type_definition_id_bnf_UNIV UNIV_I]]*)
apply (*rule order_refl*)
apply (*rule order_refl*)
apply (*rule F.ctor_rel_induct[unfolding rel_pre_F_def id_apply vimage2p_def o_apply]*)
apply (*erule mp[OF spec2[OF IH2], unfolding vimage2p_def o_apply rel_pre_GF_M_def type_definition.Abs_inverse[OF BNF_Composition.type_definition_id_bnf_UNIV UNIV_I]]*)
apply (*rule F.ctor_rel_induct[unfolding rel_pre_F_def id_apply vimage2p_def o_apply]*)
apply (*erule mp[OF spec2[OF IH2], OF vimage2p_mono[OF _ pre_GF_M.rel_mono], unfolding vimage2p_def o_apply rel_pre_GF_M_def type_definition.Abs_inverse[OF BNF_Composition.type_definition_id_bnf_UNIV UNIV_I]]*)
apply (*rule order_refl*)
apply *assumption*
apply (*rule order_refl*)
done

lemmas *n2m_ctor_induct_G_G_F = spec[OF spec [OF*
n2m_rel_induct_G_G_F[of (=) BNF_Def.Grp (Collect R) id BNF_Def.Grp (Collect S) id for R S,
unfolding G.rel_eq F.rel_eq eq_le_Grp_id_iff_all_simps(1,2)[symmetric]]],
unfolding eq_alt pre_G_M.rel_Grp pre_GF_M.rel_Grp pre_G_M.map_id0 pre_GF_M.map_id0,
unfolding vimage2p_comp vimage2p_id_comp_apply_comp_id Grp_id_mono_subst
type_copy_vimage2p_Grp_Rep[OF BNF_Composition.type_definition_id_bnf_UNIV]
type_copy_Abs_o_Rep[OF BNF_Composition.type_definition_id_bnf_UNIV]
eqTrueI[OF subset_UNIV] simp_thms(22)
atomize_conjL[symmetric] atomize_all[symmetric] atomize_imp[symmetric],
unfolding subset_iff mem_Collect_eq]

8 Mutual View on Nested Coatypes

bnf-axiomatization ('a, 'b) *coF0*
bnf-axiomatization ('a, 'b) *coG0*

8.1 Nested definition

codatatype 'a *coF* = *CcoF ('a, 'a coF) coF0*
codatatype 'a *coG* = *CcoG ('a, ('a coG) coF) coG0*

type-synonym ('b, 'c) *coF_pre_coF* = ('c, 'b) *coF0*
type-synonym ('c, 'a) *coG_pre_coG* = ('a, 'c *coF*) *coG0*

term *dtor_unfold_coF* :: ('b ⇒ ('b, 'c) *coF_pre_coF*) ⇒ 'b ⇒ 'c *coF*
term *dtor_unfold_coG* :: ('c ⇒ ('c, 'a) *coG_pre_coG*) ⇒ 'c ⇒ 'a *coG*
term *dtor_corec_coF* :: ('b ⇒ ('c *coF* + 'b, 'c) *coF_pre_coF*) ⇒ 'b ⇒ 'c *coF*
term *dtor_corec_coG* :: ('c ⇒ ('a *coG* + 'c, 'a) *coG_pre_coG*) ⇒ 'c ⇒ 'a *coG*
thm *coF.dtor_rel_coinduct*
thm *coG.dtor_rel_coinduct[unfolded rel_pre_coG_def id_apply]*

8.2 Isomorphic Mutual Definition

codatatype 'a *coG_M* = *CcoG* ('a, 'a *coGcoF_M*) *coG0*
and 'a *coGcoF_M* = *CcoF* ('a *coG_M*, 'a *coGcoF_M*) *coF0*

type-synonym ('b, 'c) *coGcoF_M_pre_coGcoF_M* = ('c, 'b) *coF0*
type-synonym ('c, 'a) *coG_M_pre_coG_M* = ('a, 'c) *coG0*

term *dtor_unfold_coG_M* :: ('b ⇒ ('c, 'a) *coG_M_pre_coG_M*) ⇒ ('c ⇒ ('c, 'b) *coGcoF_M_pre_coGcoF_M*) ⇒ 'b ⇒ 'a *coG_M*
term *dtor_unfold_coGcoF_M* :: ('b ⇒ ('c, 'a) *coG_M_pre_coG_M*) ⇒ ('c ⇒ ('c, 'b) *coGcoF_M_pre_coGcoF_M*) ⇒ 'c ⇒ 'a *coGcoF_M*
term *dtor_corec_coG_M* :: ('b ⇒ ('a *coGcoF_M* + 'c, 'a) *coG_M_pre_coG_M*) ⇒ ('c ⇒ ('a *coGcoF_M* + 'c, 'a *coG_M* + 'b) *coGcoF_M_pre_coGcoF_M*) ⇒ 'b ⇒ 'a *coG_M*
term *dtor_corec_coGcoF_M* :: ('b ⇒ ('a *coGcoF_M* + 'c, 'a) *coG_M_pre_coG_M*) ⇒ ('c ⇒ ('a *coGcoF_M* + 'c, 'a *coG_M* + 'b) *coGcoF_M_pre_coGcoF_M*) ⇒ 'c ⇒ 'a *coGcoF_M*
thm *coG_M_coGcoF_M.dtor_rel_coinduct[unfolded rel_pre_coG_M_def rel_pre_coGcoF_M_def]*

8.3 Mutualization

8.3.1 Coiterators

definition *n2m_dtor_unfold_coG* :: ('b ⇒ ('c, 'a) *coG_M_pre_coG_M*) ⇒ ('c ⇒ ('c, 'b) *coGcoF_M_pre_coGcoF_M*) ⇒ 'b ⇒ 'a *coG*

where *n2m_dtor_unfold_coG s1 s2* = *dtor_unfold_coG* (*BNF_Composition.id_bnf* o *BNF_Composition.id_bnf* o *map_pre_coG_M id* (*id* :: *unit* ⇒ *unit*) (*dtor_unfold_coF* (*BNF_Composition.id_bnf* o *BNF_Composition.id_bnf* o *s2*)) o *s1*)

definition *n2m_dtor_unfold_coG_coF* :: ('b ⇒ ('c, 'a) *coG_M_pre_coG_M*) ⇒ ('c ⇒ ('c, 'b) *coGcoF_M_pre_coGcoF_M*) ⇒ 'c ⇒ 'a *coG coF*

where *n2m_dtor_unfold_coG_coF s1 s2* = *dtor_unfold_coF* (*BNF_Composition.id_bnf* o *BNF_Composition.id_bnf* o *map_pre_coGcoF_M id* (*id* :: *unit* ⇒ *unit*) (*n2m_dtor_unfold_coG s1 s2*) *id* o *s2*)

lemma *coG_dtor_o_unfold*: *dtor_coG* o *dtor_unfold_coG* *s* = *map_pre_coG id* (*dtor_unfold_coG* *s*) o *s*
unfolding *fun_eq_iff o_apply coG.dtor_unfold* **by** *simp*

lemma *coF_dtor_o_unfold*: *dtor_coF* o *dtor_unfold_coF* *s* = *map_pre_coF id* (*dtor_unfold_coF* *s*) o *s*
unfolding *fun_eq_iff o_apply coF.dtor_unfold* **by** *simp*

lemma *coG_dtor_o_corec*: *dtor_coG* o *dtor_corec_coG* *s* = *map_pre_coG id* (*case_sum id* (*dtor_corec_coG* *s*)) o *s*

unfolding *fun_eq_iff o_apply coG.dtor_corec* **by** *simp*

lemma *coF_dtor_o_corec*: *dtor_coF* o *dtor_corec_coF* *s* = *map_pre_coF id* (*case_sum id* (*dtor_corec_coF* *s*)) o *s*

unfolding *fun_eq_iff o_apply coF.dtor_corec* **by** *simp*

lemma *n2m_dtor_unfold_coG*:

dtor_coG o *n2m_dtor_unfold_coG s1 s2* = *BNF_Composition.id_bnf* o *BNF_Composition.id_bnf* o *map_pre_coG_M id* (*id* :: *unit* ⇒ *unit*) (*dtor_unfold_coF* (*BNF_Composition.id_bnf* o *BNF_Composition.id_bnf* o *s2*)) o *s1*

unfolding *n2m_dtor_unfold_coG_def n2m_dtor_unfold_coG_coF_def*
map_pre_coG_M_def map_pre_coF_M_def map_pre_coG_M_def map_pre_coGcoF_M_def
coG_dtor_o_unfold id_apply comp_id id_comp comp_assoc

rewriteL_comp_comp[OF type_copy_map_comp0_undo[OF BNF_Composition.type_definition_id_bnf_UNIV BNF_Composition.type_definition_id_bnf_UNIV BNF_Composition.type_definition_id_bnf_UNIV pre_coG_M.map_comp0[unfolding map_pre_coG_M_def]]]

$coF.dtor_unfold_o_map$
 $rewriteL_comp_comp[OF type_copy_Rep_o_Abs[OF BNF_Composition.type_definition_id_bnf_UNIV]] ..$

lemma $n2m_dtor_unfold_coG_coF$:

$dtor_coF o n2m_dtor_unfold_coG_coF s1 s2 = BNF_Composition.id_bnf o BNF_Composition.id_bnf o map_pre_coGcoF_M$
 $id (n2m_dtor_unfold_coG s1 s2) (n2m_dtor_unfold_coG_coF s1 s2) o s2$

unfolding $n2m_dtor_unfold_coG_coF_def map_pre_coF_def map_pre_coG_M_def map_pre_coGcoF_M_def$
 $coF_dtor_o_unfold id apply comp_id id comp comp_assoc$
 $rewriteL_comp_comp[OF coF0.map_comp0[symmetric]]$
 $rewriteL_comp_comp[OF type_copy_Rep_o_Abs[OF BNF_Composition.type_definition_id_bnf_UNIV]] ..$

8.3.2 Corecursors

definition $n2m_dtor_corec_coG$::

$('b \Rightarrow ('a coG coF + 'c, 'a) coG_M_pre_coG_M) \Rightarrow ('c \Rightarrow ('a coG coF + 'c, 'a coG + 'b) coGcoF_M_pre_coGcoF_M)$
 $\Rightarrow 'b \Rightarrow 'a coG$

where $n2m_dtor_corec_coG s1 s2 =$

$dtor_corec_coG (BNF_Composition.id_bnf o BNF_Composition.id_bnf o$
 $map_pre_coG_M id (id :: unit \Rightarrow unit))$
 $(case_sum (map_coF Inl) (dtor_corec_coF (BNF_Composition.id_bnf o BNF_Composition.id_bnf o$
 $map_pre_coGcoF_M (id :: unit \Rightarrow unit) id (map_sum (map_coF Inl) id) o s2))) o$
 $s1)$

definition $n2m_dtor_corec_coG_coF$::

$('b \Rightarrow ('a coG coF + 'c, 'a) coG_M_pre_coG_M) \Rightarrow ('c \Rightarrow ('a coG coF + 'c, 'a coG + 'b) coGcoF_M_pre_coGcoF_M)$
 $\Rightarrow 'c \Rightarrow 'a coG coF$

where $n2m_dtor_corec_coG_coF s1 s2 = dtor_corec_coF (BNF_Composition.id_bnf o BNF_Composition.id_bnf$
 $o map_pre_coGcoF_M (id :: unit \Rightarrow unit) (case_sum id (n2m_dtor_corec_coG s1 s2))) id o s2)$

lemma $n2m_dtor_corec_coG$:

$dtor_coG o n2m_dtor_corec_coG s1 s2 = BNF_Composition.id_bnf o BNF_Composition.id_bnf o map_pre_coG_M$
 $id (case_sum id (n2m_dtor_corec_coG_coF s1 s2)) o s1$

unfolding $n2m_dtor_corec_coG_def n2m_dtor_corec_coG_coF_def$
 $map_pre_coG_def map_pre_coF_def map_pre_coG_M_def map_pre_coGcoF_M_def$
 $coG_dtor_o_corec$
 $id apply comp_id id comp comp_assoc[symmetric] map_sum.comp map_sum.id$
 $case_sum_o_inj(1) case_sum_o_map_sum o case_sum$
 $rewriteR_comp_comp[OF coG0.map_comp0[symmetric]]$
 $rewriteR_comp_comp[OF coF0.map_comp0[symmetric]]$
 $coF.map_comp0[symmetric] coF.map_id0$
 $coF.dtor_corec_o_map$
 $rewriteR_comp_comp[OF type_copy_Rep_o_Abs[OF BNF_Composition.type_definition_id_bnf_UNIV]] ..$

lemma $n2m_dtor_corec_coG_coF$:

$dtor_coF o n2m_dtor_corec_coG_coF s1 s2 = BNF_Composition.id_bnf o BNF_Composition.id_bnf o map_pre_coGcoF_M$
 $id (case_sum id (n2m_dtor_corec_coG s1 s2)) (case_sum id (n2m_dtor_corec_coG_coF s1 s2)) o s2$

unfolding $n2m_dtor_corec_coG_coF_def map_pre_coF_def map_pre_coG_M_def map_pre_coGcoF_M_def$
 $coF_dtor_o_corec id apply comp_id id comp comp_assoc$
 $rewriteL_comp_comp[OF coF0.map_comp0[symmetric]]$
 $rewriteL_comp_comp[OF type_copy_Rep_o_Abs[OF BNF_Composition.type_definition_id_bnf_UNIV]] ..$

8.3.3 Coinduction

lemma $n2m_rel_coinduct_coG_coG_coF$:

assumes $CIH1: \forall x y. R x y \longrightarrow BNF_Def.vimage2p (BNF_Composition.id_bnf o BNF_Composition.id_bnf)$
 $(BNF_Composition.id_bnf o BNF_Composition.id_bnf) (rel_pre_coG_M P R S) (dtor_coG x) (dtor_coG y)$

and $CIH2: \forall x y. S x y \longrightarrow BNF_Def.vimage2p (BNF_Composition.id_bnf o BNF_Composition.id_bnf)$
 $(BNF_Composition.id_bnf o BNF_Composition.id_bnf) (rel_pre_coGcoF_M P R S) (dtor_coF x) (dtor_coF y)$

shows $R \leq rel_coG P \wedge S \leq rel_coF (rel_coG P)$

apply (rule context_conjI)

apply (rule coG.dtor_rel_coinduct[unfolded rel_pre_coG_def id_apply vimage2p_def o_apply])

apply (erule mp[OF spec2[OF CIH1], THEN vimage2p_mono[OF _pre_coG_M.rel_mono], unfolded vimage2p_def

$o_apply rel_pre_coG_M_def type_definition.Abs_inverse[OF BNF_Composition.type_definition_id_bnf_UNIV UNIV_I]])$

apply (rule order_refl)

```

apply (rule order_refl)
apply (rule coF.dtor_rel_coinduct[unfolded rel_pre_coF_def id_apply vimage2p_def o_apply])
apply (erule mp[OF spec2[OF CIH2], unfolded vimage2p_def o_apply rel_pre_coGcoF_M_def type_definition.Abs_inverse[OF
BNF_Composition.type_definition_id_bnf_UNIV_UNIV_I]])

```

```

apply (rule coF.dtor_rel_coinduct[unfolded rel_pre_coF_def id_apply vimage2p_def o_apply])
apply (erule mp[OF spec2[OF CIH2], THEN vimage2p_mono[OF _ pre_coGcoF_M.rel_mono], unfolded vimage2p_def o_apply rel_pre_coGcoF_M_def type_definition.Abs_inverse[OF BNF_Composition.type_definition_id_bnf_UNIV
UNIV_I]])
apply (rule order_refl)
apply assumption
apply (rule order_refl)
done

```

```

lemmas n2m_ctor_induct_coG_coG_coF = spec[OF spec[OF spec[OF spec[OF
n2m_rel_coinduct_coG_coG_coF[of _ (=),
unfolded coG.rel_eq coF.rel_eq le_fun_def le_bool_def all_simps(1,2)[symmetric]]]]]]]]

```